

Math Basics

Digit
Concept

int n = 7789 % 10 = 9
 \downarrow
 110 778.9 % 10 = 8
 \downarrow
 110 77 % 10 = 7
 \downarrow
 110 7 % 10 = 7
 \downarrow
 110 0

extraction of digits is taking input 'n' and while $n > 0$, and prints the last digit % 10. Then divides n by 10 to iterate till 0. *IN REVERSE ORDER

Count
Digits

```
count = 0;
while (n > 0) {
    lastDigit = n % 10;
    count += 1;
    n = n / 10;
}
```

* just make a counter for the coded loop above

* time complexity is $O(\log_{10}(N))$ b/c loop is getting run number of times it is divided by 10

* If division is happening by 10, $\log_{10}(N)$. by 2, $\log_2(N)$

Reverse
a
Number

n = 7789 → 9877

```
reverseNum = 0;
lastDigit = n % 10;
```

Digit
Concept

```
reverseNum = (reverseNum * 10) + lastDigit;
}
return reverseNum;
```


Day 9

2/19/25

Check
Palindrome

palindrome is ex $121 == 121$, $12321 == 12321$

take reverse number and compare w/ original

Armstrong
Number

$$n = 371 = 3^3 + 7^3 + 1^3 = 371$$

make sum = 0 variable

sum = sum + (lastDigit**3);

run rest of code the same

Print
All Divisors
 $O(N)$

36 \rightarrow 1, 2, 3, 4, 6, 9, 12, 18, 36 // divide w/o remainder

loop from $1 \rightarrow n$ for($i=1; i \leq n, i++$) {
check if leaves 0 remainder if($n \% i == 0$) print i;
}

first check if i is factor

also check if the factor is not equal to i , so no duplicate

Sorting

store in a list vector<int> ls;
sort list for(auto it: ls) cout << it << " ";
tc is $O(\# \text{ of factors} \rightarrow \log n)$

Check
For
Prime

$O(N)$
Brute
Force

*check if factors are 1 & itself

int cnt = 0

for($i=1; i \leq n, i++$) {

if($n \% i == 0$) cnt++;

if(cnt == 2) ✓;

else X;

if($n \% i == 0$)

cnt++;

if($(n/i) !=$

cnt++;

if(cnt == 2) ✓;

else X;

2/20/25

 $N1 = 9$ $N2 = 12$ ↓
1, 3, 9↓
1, 2, 6, 12, 3, 4

highest divisor

so $\text{gcd}(9, 12) = 3$

From $1 \rightarrow n2$, check every number to see if it divides both, keep replacing gcd till highest number

```
for (i = 1; i <= min(n1, n2); i++) {
    if (n1 % i == 0 && n2 % i == 0) {
        gcd = i;
    }
}
```

} }
flipped ↓

 $TC = O(\min(n1, n2))$

```
for (i = min(n1, n2); i >= 1; i--) {
    if (n1 % i == 0 && n2 % i == 0) {
        print i;
        break; // same time complexity
    }
}
```

} }

EQUILATERAL ALGORITHM

$$\text{gcd}(a, b) = \text{gcd}(a - b, b)$$

$$\text{gcd}(10, 5) = \text{gcd}(5, 5) \rightarrow \text{gcd}(5 - 5, 5) = (0, 5)$$

so when $a > b$

$$\text{gcd}(a, b) = \text{gcd}(a \% b, b)$$

* if one value = 0, the other value is the GCF