



Adventures & Explorations: LC

Guideline

- <https://www.teamblind.com/post/New-Year-Gift---Curated-List-of-Top-75-LeetCode-Questions-to-Save-Your-Time-OaM1orEU>
- Steps:
 - Keep calm and patient
 - Read question carefully and understand requirements fully: 5-10 mins
 - Set out strategy and get confirmation
 - Code main logic and handle edge cases (don't hope for luck)
- Good advice/watching: <https://www.youtube.com/c/NeetCode> <https://neetcode.io/>
- udemy course:
<https://www.udemy.com/course/datastructurescncpp/?referralCode=BD2EF8E61A98AB5E011D>

Key steps:

- #1: Read Q clearly and understand requirements exactly, ask for clarification as needed, figure which category and which what data structure.
- #2: Discuss and confirm strategy
- #3: Code main logic first, and try to cover edge cases
- #4: don't hope for luck that it works magically, make sure/ understand every step clearly.
- #5: Correct/Improve as it runs test cases

C++ STL: unordered_map

```
#include <unordered_map>
unordered_map<string, int> umap;
unordered_map<string, int>::iterator itr;
pair<string, int> elem;
umap[“test1”] = 1; umap[“test2”] = 2;
elem = make_pair(“test3”, 199); umap.insert(elem); // add

itr = umap.find(“mykey”); if (itr != umap.end()) printf(“found mykey”); //look up

for (itr =umap.begin(); itr != umap.end(); itr++) { // iterate thru for (auto kv: umap) kv.first , kv.second
    printf(“key is %s , val is %d \n”, itr->first, itr->second);
umap.erase(“my key”); umap.erase(umap.begin());
Umap.count(“my key”); return zero if not found;
umap.size(); // return number of elements in the map
Umap.empty() ;//tell if there is anything
```

C++ STL unordered_set

```
#include <unordered_set>
unordered_set<string> set1;
unordered_set<string>::iterator itr1;
unordered_set<int> set2;
set1.insert("hello"); set1.insert("world"); // add
if (set1.find("myname") != set1.end()) { // found it ... } // look up
for (auto itr = set1.begin(); itr != set1.end(); itr++) { //iterate
    printf("%s", *itr);
}
set2.insert(10); set2.insert(20); set2.erase(20); set2.erase(set2.find(10));
set2.count(key) => 0 or 1; set2.size() how many elements?
if (set1.empty()) { // nothing inside ...}
```

C++ STL Stack and Queue

```
#include <stack>
stack<int> st;
st.push(10); st.push(20); st.push(30);
while (!st.empty()) { printf("%d", st.top()); st.pop();}
st.size();
```

```
#include <queue>
queue<int> q;
q.push(100); q.push(200); q.push(300);
while (!q.empty()) { printf("%d, ", q.front()); q.pop();}
q.size();
```

C++ STL Vector: used as ArrayList in Java

```
#include <vector>
vector<int> g1;
for (int i = 1; i <= 10; i++) g1.push_back(i * 10);
cout << "\n Reference operator [g] : g1[2] = " << g1[2];
cout << "\n Using at : g1.at(4) = " << g1.at(4);
cout << "\nfront() : g1.front() = " << g1.front();
cout << "\nback() : g1.back() = " << g1.back();
int* pos = g1.data(); // pointer to the first element

g1.push_back(15); g1.pop_back();
g1.insert(g1.begin(), 100); g1.insert(g1.begin()+3, 300);
g1.erase(g1.begin() +3); //erase
int index = find(g1.begin(), g1.end(), key) - g1.begin();
#include <bits/stdc++.h>
vector<int> v1 { 1, 20, 3, 40, 5, 60};
sort(v1.begin(), v1.end(), less<int>()); // less is default { 1,3,5,20, 40, 60}
```

C++ STL string class

```
#include <string>
std::string s = "Hello";
std::string greet = s + " World"; //concatenation easy!
str.push_back('s'); str.length(); str.at(i) is same as str[i]
to_string(123); // converte integer to string
l = stoi("123"); // return interger 123;
str.find(subStr,0) == 0 => str is started with subStr
str.push_back('c');
string::append (size_type num, char c)

str.substr(pos, len);
const char *cp = str.c_str();
```

C++ STD priority_queue: min heap, adjust on the fly based on new value pq.push(), pq.pop()

- #include <queue>
- priority_queue<int, vector<int>, Cmp> pq;
// without Cmp, big num# first pop(), Cmp=less<int>, called Max heap
- Comparator:

```
class Cmp {  
public:  
    bool operator()(int i1, int i2) {  
        return i1 > i2 ; //pq.top(): will be smallest number, min heap  
    }  
}
```
- priority_queue<int, vector<int>, greater<int>> min_hep;

Binary search in vector

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <iostream>

typedef std::vector<int>::iterator iter;

int main() {

    std::vector<int> vec = {10, 20, 30, 30, 20, 10, 10, 20};

    // sort the data
    // the data will be:
    // 10, 10, 10, 20, 20, 20, 30, 30
    std::sort(vec.begin(), vec.end());

    // index of the first element, greater than or equal to 20
    iter low = std::lower_bound(vec.begin(), vec.end(), 20);      Return 3

    // index of the first element, greater than 20
    iter high = std::upper_bound(vec.begin(), vec.end(), 20);      Return 6

    std::cout << "index of first element, greater than or equal to 20 is: " << (low - vec.begin()) << '\n';

    std::cout << "index of first element, greater than to 20 is: " << (high - vec.begin()) << '\n';

    // classic binary search
    // check whether a given value exists in the array or not
    if (std::binary_search(vec.begin(), vec.end(), 99)) {
        std::cout << "Found\n";
    } else {
        std::cout << "Not found\n";
    }
}
```

- sort: you can use binary search only on a *sorted* data, so you must guarantee that the data is sorted, before searching.
- lower_bound: this function returns an *iterator* to the first element that is **greater than or equal to** value.
- upper_bound: this function returns an *iterator* to the first element that is **greater than** value.
- binary_search: this function returns a *boolean*, whether the value is found or not (exactly as your program).

Strings

```
char str[] = "HELLO";
```



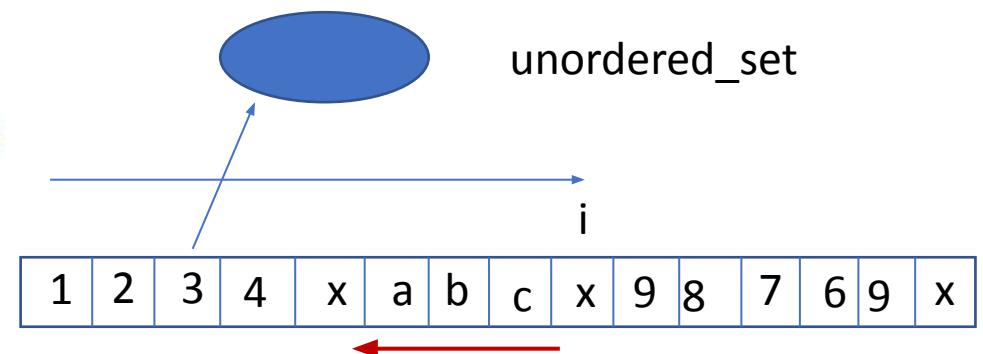
Beginning Of String

End Of String

LC String#1: 3. Longest Substring Without Repeating Characters

```
1 class Solution {
2     public:
3         int lengthOfLongestSubstring(string s) {
4             int res = 0;
5             int c = 0;
6             if (s.length() < 2) return s.length();
7             unordered_set<char> seen;
8             for (int i = 0; i < s.length(); i++) {
9                 if (seen.find(s.at(i)) == seen.end()) {
10                     c++;
11                     res = max(c, res);
12                     seen.insert(s.at(i));
13                 } else {
14                     seen.clear();
15                     // search back for the last occurrence of this character
16                     seen.insert(s.at(i));
17                     c = 1;
18                     for (int j = i-1; j >= 0; j--) {
19                         if (s.at(j) != s.at(i)) {
20                             seen.insert(s.at(j));
21                             c++;
22                         } else {
23                             break;
24                         }
25                     }
26                 }
27             }
28             return res;
29         }
30     };
```

- Scan thru string, push each one into “set” named as “seen” if not seen, c++
- If it is seen before, reset count “c”, and clear “seen”, search back until current char is hit(**where is i is adjusted to**) add all back into “seen”



LC String#2: two pointers/ Caterpillar algorithm

424. Longest Repeating Character with K Replacement

```
1 class Solution {
2 public:
3     int characterReplacement(string s, int k) {
4         int L = s.length();
5         int res = 0;
6         int l = 0, r = 0;
7         int counters[26] = { 0 };
8         bool adv_right = true;
9         while (l <= r && r < L) {
10             if (adv_right)
11                 counters[s.at(r)-'A']++;
12             else
13                 counters[s.at(l-1)-'A']--;
14
15             //find max freq of current window
16             int max_freq = 0;
17             for (int i = 0; i < 26; i++) {
18                 max_freq = max(max_freq, counters[i]);
19             }
20             int rc = r-l+1 - max_freq;
21             if (rc <= k) {
22                 res = max(r-l+1, res);
23                 printf("\n res=%d, l=%d, r=%d", res, l,r);
24                 adv_right = true;
25                 r++;
26             } else {
27                 printf("\n advance left l=%d", l+1);
28                 l++;
29                 adv_right = false;
30             }
31         }
32     return res;
33 }
34 }
```

You are given a string `s` and an integer `k`. You can choose any character of the string and change it to any other uppercase English character. You can perform this operation at most `k` times.

Return the length of the longest substring containing the same letter you can get after performing the above operations.

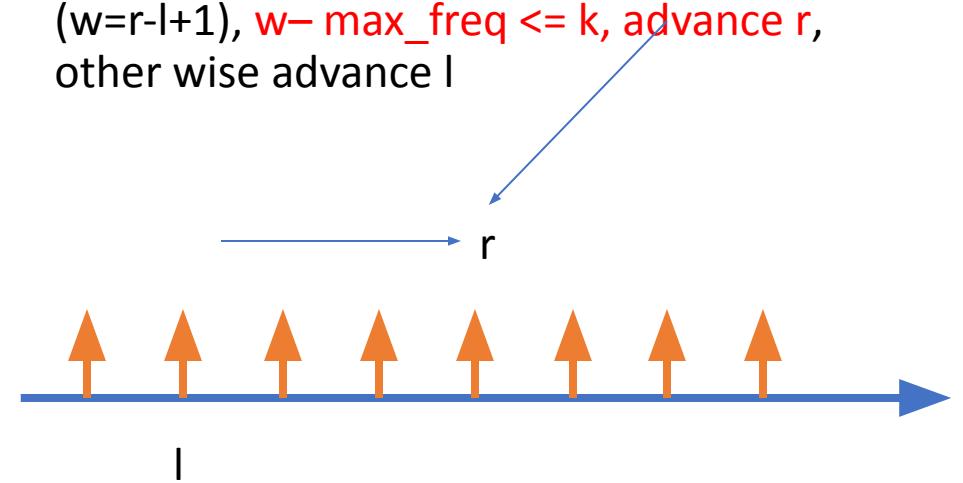
Example 1:

Input: `s = "ABAB"`, `k = 2`

Output: 4

Explanation: Replace the two 'A's with two 'B's or vice versa.

- Use Sliding windows: $l=0, r=0 \rightarrow L$
- Use a `counters[26]` array to count each **inside window's letters** frequency when l/r is advanced
- Inside a sliding window, find the `max_freq` letter and its counter, window width ($w=r-l+1$), **$w - max_freq \leq k$, advance r** , otherwise advance l



LC string#3: two pointers/ Caterpillar algorithm

76. Minimum Window Substring: from S including T

```
1 class Solution {
2     public:
3         string minWindow(string s, string t) {
4             int l = 0, r = 0, res= INT_MAX, r_ = 0, l_ = 0;
5             bool adv_r = true;
6             int LEN = 'z' - 'A' + 1;
7             int t_counters['z' - 'A' + 1] = { 0 };
8             int s_counters['z' - 'A' + 1] = { 0 };
9
10            for (int i = 0; i < t.length(); i++) t_counters[t.at(i)-'A']++;
11            while (l <= r && r < s.length()) {
12                if (adv_r)
13                    s_counters[s.at(r)-'A']++;
14                else
15                    s_counters[s.at(l-1)-'A']--;
16
17                //check if current window has all string t
18                bool cover_t = true;
19                for (int i = 0; i < LEN; i++) {
20                    if (t_counters[i] != 0 && s_counters[i] < t_counters[i]) {
21                        cover_t = false;
22                        break;
23                    }
24                }
25                if (cover_t) {
26                    if (r-l+1 < res) {
27                        r_ = r; l_ = l;
28                        res = r - l + 1;
29                    }
30                    adv_r = false;
31                    l++;
32                } else {
33                    adv_r = true;
34                    r++;
35                }
36            }
37
38            return res == INT_MAX? "":s.substr(l_, res);
39        }
40    };
```

Given two strings s and t of lengths m and n respectively, return the **minimum window substring** of s such that every character in t (**including duplicates**) is included in the window. If there is no such substring, return the empty string "".

The testcases will be generated such that the answer is **unique**.

A **substring** is a contiguous sequence of characters within the string.

Example 1:

Input: $s = \text{"ADOBECODEBANC"}$, $t = \text{"ABC"}$

Output: "BANC"

Explanation: The minimum window substring "BANC" includes 'A', 'B', and 'C' from string t .

- **Sliding window:** expand R to cover T, write down W
- Advance L until not cover T, write down smaller W,
- Go back to first step until reach the end of string.

LC string#4: 242. Valid Anagram: s = "anagram", t = "nagaram" => true

```
1 * class Solution {
2 public:
3     bool isAnagram(string s, string t) {
4         if (s.length() != t.length()) return false;
5         int s_c['z'-'a'+1] = {0};
6         int t_c['z'-'a'+1] = {0};
7         for (int i = 0; i < s.length(); i++) {
8             s_c[s.at(i)-'a']++;
9             t_c[t.at(i)-'a']++;
10        }
11        return memcmp(s_c, t_c, sizeof(int)*('z'-'a'+1)) == 0;
12    }
13}
```

An Anagram is a word or phrase formed by rearranging the letters of a different word or phrase, typically using all the original letters exactly once.

LC string#5: 49. Group Anagrams:

```
1 class Solution {
2 public:
3     vector<vector<string>> groupAnagrams(vector<string>& strs) {
4         vector<vector<string>> res;
5         vector<vector<int>> counters;
6         vector<bool> visited;
7         for (int i = 0; i < strs.size(); i++) {
8             string s = strs[i];
9             vector<int> c(26,0);
10            for (int j = 0; j < s.length(); j++)
11                c[s.at(j) - 'a']++;
12            counters.push_back(c);
13            visited.push_back(false);
14        }
15
16        for (int i = 0; i < strs.size(); i++) {
17            if (!visited[i]) {
18                vector<string> ans;
19                ans.push_back(strs[i]);
20                visited[i] = true;
21                for (int j = i+1; j < strs.size(); j++) {
22                    if (counters[i] == counters[j]) {
23                        ans.push_back(strs[j]);
24                        visited[j] = true;
25                    }
26                }
27                res.push_back(ans);
28            }
29        }
30    }
31    return res;
32 };
```

```
Input: strs = ["eat","tea","tan","ate","nat","bat"]
Output: [[["bat"]], [ ["nat", "tan"] ], [ ["ate", "eat", "tea"] ]]
```

- **Character counters**: for each string **vector<int> c(26,0)**
- Visited[] Boolean to speed up
- Scan thru strs and counters[]
vector compare: **counters[i] == counters[j]**

LC string#6: 20. Valid Parentheses:

```
1 class Solution {
2 public:
3     bool isValid(string s) {
4         stack<char> st;
5
6         for (int i = 0; i < s.length(); i++) {
7             if (s.at(i) == '(' || s.at(i) == '[' || s.at(i) == '{')
8                 st.push(s.at(i));
9             else {
10                 if (st.empty()) return false;
11                 char c = st.top(); st.pop();
12                 if ((c == '{' && s.at(i) == '}') ||
13                     (c == '[' && s.at(i) == ']') ||
14                     (c == '(' && s.at(i) == ')'))
15                 ) {
16                     continue;
17                 } else
18                     return false;
19             }
20         }
21         return st.empty();
22     }
23 }
```

Input: s = "()[]{}"
Output: true

Input: s = "([)]"
Output: false

- Use stack: opening one, push, closing one pop and compare.

LC string#7: 125. Valid Palindrome:

```
1 class Solution {
2     public:
3         bool isAlphaNum(char c) {
4             return ((c >= 'A' && c <= 'Z') ||
5                     (c >= 'a' && c <= 'z') ||
6                     (c >= '0' && c <= '9'));
7         }
8         bool equalIgnoreCase(char a, char b) {
9             // converte to uppercase to compare
10            if (a >= 'a' && a <= 'z') {
11                a = 'A' + (a-'a');
12            }
13            if (b >= 'a' && b <= 'z') {
14                b = 'A' + (b-'a');
15            }
16            return a == b;
17        }
18        bool isPalindrome(string s) {
19            int l = 0;
20            int r = s.length() -1;
21
22            while (l < r) {
23                while (l < s.length() && !isAlphaNum(s.at(l))) l++;
24                while (r >= 0 && !isAlphaNum(s.at(r))) r--;
25                if (l >= s.length() || r < 0 ) break;
26                if (!equalIgnoreCase(s.at(l), s.at(r))) return false;
27                l++;
28                r--;
29            }
30            return true;
31        }
32    };
```

Input: s = "A man, a plan, a canal: Panama"

Output: true

Explanation: "amanaplanacanalpanama" is a palindrome.

LC string#8: 5. Longest Palindromic

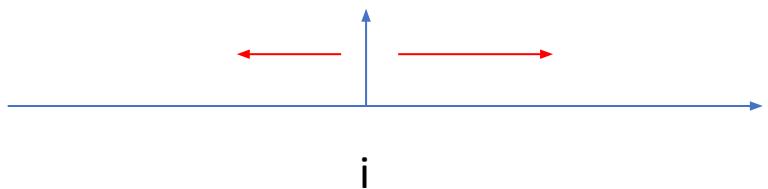
substring:

```
1 * class Solution {
2 public:
3 *     int expandHelper(string s, int l, int r) {
4 *         while (l >= 0 && r <= s.length()-1 && s.at(l) == s.at(r) ) {
5 *             l--; r++;
6 *         }
7 *         return r-l-1;
8 *     }
9 *     string longestPalindrome(string s) {
10 *         int len = 1;
11 *         int pos = 0;
12 *         for (int i = 0; i < s.length(); i++) {
13 *             int len_odd = expandHelper(s, i, i);
14 *             int len_even = expandHelper(s, i, i+1);
15 *             int max_len = max(len_odd, len_even);
16 *             if (max_len > len) {
17 *                 len = max_len;
18 *                 pos = i - (len-1)/2;
19 *             }
20 *         }
21 *         return s.substr(pos, len);
22 *     }
23 * }
```

Expand with l as center “odd”

Expand with i and i+1 as “even”

take longer one , record it along the way
to get len, pos = $i - (len-1)/2$



LC string#9: 647: Palindromic substring:

```
1 class Solution {
2     public:
3         void expandHelper(string s, int l, int r, int & result) {
4             while (l >= 0 && r < s.length() && s.at(l) == s.at(r)) {
5                 result++;
6                 l--;
7                 r++;
8             }
9         }
10        int countSubstrings(string s) {
11            int res = 0;
12
13            for (int i = 0; i < s.length(); i++) {
14                expandHelper(s, i, i, res);
15                expandHelper(s, i, i+1, res);
16            }
17
18            return res;
19        }
};
```

Given a string s , return the number of **palindromic substrings** in it.

A string is a **palindrome** when it reads the same backward as forward.

A **substring** is a contiguous sequence of characters within the string.

Example 1:

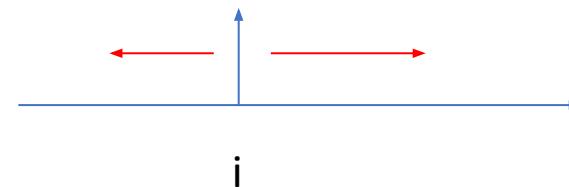
Input: $s = \text{"abc"}$

Output: 3

Explanation: Three palindromic strings: "a", "b", "c".

Expand with i as center “odd”

Expand with i and $i+1$ as “even”



LC string#10: 271: Encode and Decode strings:

```
1  class Codec {
2  public:
3
4      // Encodes a list of strings to a single string.
5      string encode(vector<string>& strs) {
6          string res;
7          for (int i = 0; i < strs.size(); i++) {
8              string tmp = to_string(strs[i].length())+"@"+strs[i];
9              res.append(tmp);
10         }
11         return res;
12     }
13
14     // Decodes a single string to a list of strings.
15     vector<string> decode(string s) {
16         string dec_s = s;
17
18         vector<string> res;
19         int i = 0;
20         while (i < dec_s.length()) {
21             int l = 0;
22             while (dec_s.at(i+l) != '@') {
23                 l++;
24             }
25             string num_s = dec_s.substr(i,l);
26             int len = stoi(num_s);
27             i += l+1; // skip @
28             res.push_back(dec_s.substr(i, len));
29             i += len;
30         }
31     }
32 }
33 }
```

Design an algorithm to encode a list of strings to a string. The encoded string is then sent over the network and is decoded back to the original list of strings.

Machine 1 (sender) has the function:

```
string encode(vector<string> strs) {
    // ... your code
    return encoded_string;
}
```

Machine 2 (receiver) has the function:

```
vector<string> decode(string s) {
    //... your code
    return strs;
}
```

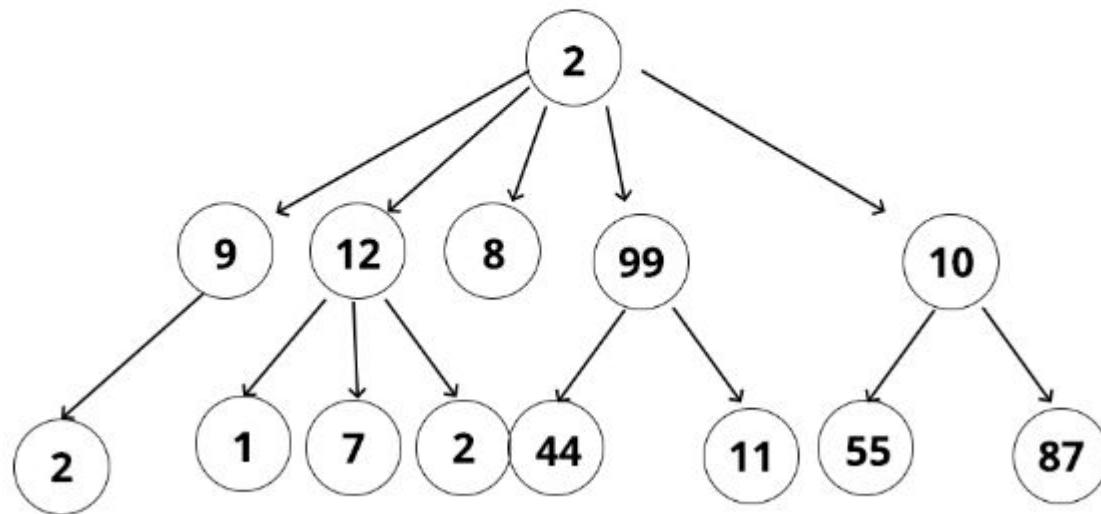
LC string#11: 13. Roman to Integer:

```
1 class Solution {
2 public:
3     int romanToInt(string s) {
4         unordered_map<string, int> map;
5
6         map["I"] = 1; map["V"] = 5; map["IV"] = 4; map["IX"] = 9;
7         map["X"] = 10; map["L"] = 50; map["XL"] = 40; map["XC"] = 90;
8         map["C"] = 100; map["D"] = 500; map["CD"] = 400; map["CM"] = 900;
9         map["M"] = 1000;
10        int i = 0;
11        int res = 0;
12        while (i < s.length()) {
13            // check if it one or two characters
14            if((i+1) < s.length()) {
15                string dStr = s.substr(i,2);
16                if (map.find(dStr) != map.end()) {
17                    res += map[dStr];
18                    i += 2;
19                    continue;
20                }
21            }
22            res += map[s.substr(i,1)];
23            i++;
24        }
25        return res;
26    }
27};
```

12. Integer to Roman

```
1 class Solution {
2 public:
3     string intToRoman(int num) {
4         vector<int> values { 1000, 900, 500, 400, 100,
5                             90, 50, 40, 10, 9, 5, 4, 1 };
6         vector<string> syms { "M", "CM", "D", "CD", "C", "XC",
7                             "L", "XL", "X", "IX", "V", "IV", "I" };
8         string res = "";
9         for (int i = 0; i < values.size() && num > 0; i++) {
10             while (num >= values[i]) {
11                 num -= values[i];
12                 res.append(syms[i]);
13             }
14         }
15         return res;
16     }
17};
```

TREES

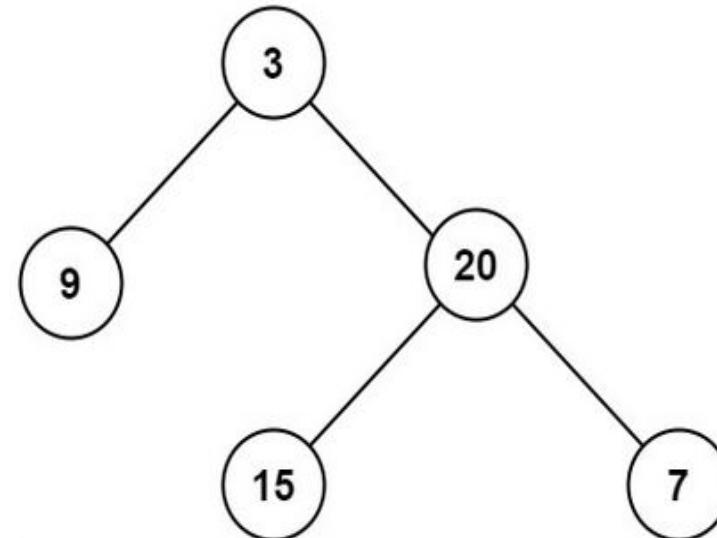


LC tree#1: 104. Maximum Depth of Binary Tree:

```
12 class Solution {  
13     public:  
14         int maxDepth(TreeNode* root) {  
15             int max_level = 0;  
16  
17             if (!root) return max_level;  
18             queue<TreeNode*> Q;  
19             Q.push(root);  
20  
21             while (!Q.empty()) {  
22                 max_level++;  
23                 int s = Q.size();  
24                 for (int i = 0; i < s; i++) {  
25                     TreeNode *n = Q.front(); Q.pop();  
26                     if (n->left) Q.push(n->left);  
27                     if (n->right) Q.push(n->right);  
28                 }  
29             }  
30  
31             return max_level;  
32         }  
33     };
```

Use Q to do a level traversal and Counting levels

Example 1:



Input: root = [3,9,20,null,null,15,7]

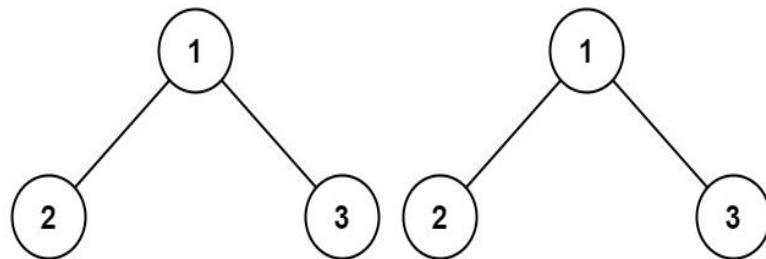
Output: 3

LC tree#2: 100. Same Tree:

```
12 class Solution {
13 public:
14     void dfs(TreeNode *root, vector<int>& v) {
15         if (root == NULL) return;
16         TreeNode *dummy = new TreeNode(INT_MAX);
17
18         stack<TreeNode*> st;
19         st.push(root);
20         while (!st.empty()) {
21             TreeNode *n = st.top(); st.pop();
22             v.push_back(n->val);
23             if (n->val != INT_MAX) {
24                 st.push(n->left? n->left:dummy);
25                 st.push(n->right? n->right:dummy);
26             }
27         }
28         delete(dummy);
29     }
30     bool isSameTree(TreeNode* p, TreeNode* q) {
31         vector<int> p_vals, q_vals;
32         dfs(p, p_vals);
33         dfs(q, q_vals);
34         return p_vals == q_vals;
35     }
36 }
```

Do a DFS using stack for each tree, for absent node, Add a dummy node with INT_MAX as flag, which should do nothing except add INT_MAX into vector after it is popped out from stack

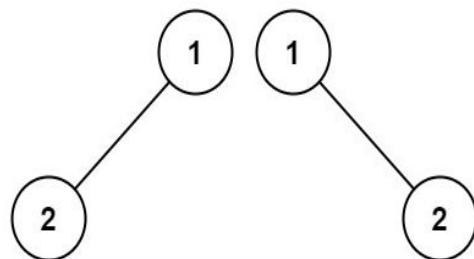
Example 1:



Input: p = [1,2,3], q = [1,2,3]

Output: true

Example 2:



Input: p = [1,2], q = [1,null,2]

Output: false

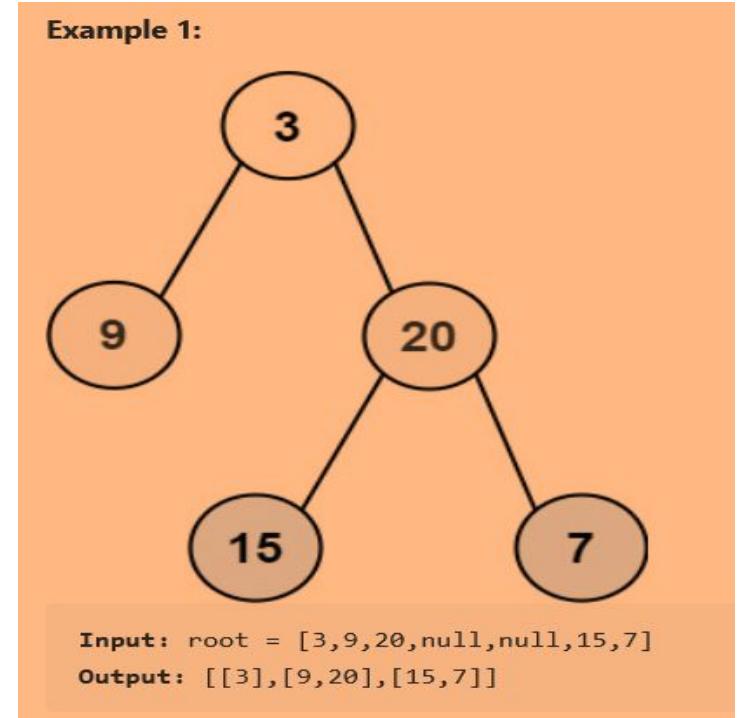
LC tree#3: 226. Invert Binary Tree:

```
12 class Solution {  
13     public:  
14         TreeNode* invertTree(TreeNode* root) {  
15             if (root == NULL) return root;  
16             TreeNode *t = root->left;  
17             root->left = root->right;  
18             root->right = t;  
19             invertTree(root->left);  
20             invertTree(root->right);  
21  
22             return root;  
23         }  
24     };
```

Exchange left with right node
Do recursion for left and right.

LC tree#4.1: 102. Binary Tree Level Order Traversal:

```
12 class Solution {
13 public:
14     vector<vector<int>> levelOrder(TreeNode* root) {
15         vector<vector<int>> res;
16         if (root == NULL) return res;
17         queue<TreeNode*> Q;
18         Q.push(root);
19
20         while (!Q.empty()) {
21             int N = Q.size();
22             vector<int> level;
23             for (int i = 0; i < N; i++) {
24                 TreeNode *n = Q.front(); Q.pop();
25                 level.push_back(n->val);
26                 if (n->left) Q.push(n->left);
27                 if (n->right) Q.push(n->right);
28             }
29             res.push_back(level);
30         }
31
32         return res;
33     }
34 }
```



Using **Q** to do level traversal
Check **Q size** in the beginning of loop,
which is number of nodes in that level.

LC string#4.2: 144. Binary Tree preOrder Traversal:

```
12 class Solution {  
13     public:  
14         vector<int> preorderTraversal(TreeNode* root) {  
15             vector<int> res;  
16             if (root == NULL) return res;  
17 #ifdef RECURSIVE  
18             res.push_back(root->val);  
19             vector<int> left = preorderTraversal(root->left);  
20             for (int i = 0; i < left.size(); i++)  
21                 res.push_back(left[i]);  
22             vector<int> right = preorderTraversal(root->right);  
23             for (int i = 0; i < right.size(); i++)  
24                 res.push_back(right[i]);  
25 #endif  
26             TreeNode *cur = root;  
27             stack<TreeNode *> st;  
28             while (cur != NULL || !st.empty()) {  
29                 if (cur != NULL) {  
30                     res.push_back(cur->val);  
31                     st.push(cur);  
32                     cur= cur->left; // push down all the way to bottom most left node  
33                 } else {  
34                     cur = st.top(); st.pop();  
35                     cur = cur->right;  
36                 }  
37             }  
38         }  
39         return res;  
40     }  
41 };
```

DFS using stack is cleanest!

```
42     TreeNode *cur = root;  
43     stack<TreeNode *> st;  
44     st.push(root);  
45     while (!st.empty()) {  
46         TreeNode *n = st.top(); st.pop();  
47         res.push_back(n->val);  
48         if (n->right) st.push(n->right);  
49         if (n->left) st.push(n->left);  
50     }  
51     return res;  
52 }
```

LC string#4.3: 94. Binary Tree Level InOrder Traversal:

```
12  class Solution {
13  public:
14      vector<int> inorderTraversal(TreeNode* root) {
15         vector<int> res;
16         if (root == NULL) return res;
17 #ifdef RECURSIVE
18         vector<int> left = preorderTraversal(root->left);
19         for (int i = 0; i < left.size(); i++)
20             res.push_back(left[i]);
21         res.push_back(root->val);
22         vector<int> right = preorderTraversal(root->right);
23         for (int i = 0; i < right.size(); i++)
24             res.push_back(right[i]);
25 #endif
26         TreeNode *cur = root;
27         stack<TreeNode *> st;
28         while (cur != NULL || !st.empty()) {
29             if (cur != NULL) {
30                 st.push(cur);
31                 cur= cur->left; // push down all the way to bottom most left node
32             } else {
33                 cur = st.top(); st.pop();
34                 res.push_back(cur->val);
35                 cur = cur->right;
36             }
37         }
38     }
39     return res;
40 }
41 };
```

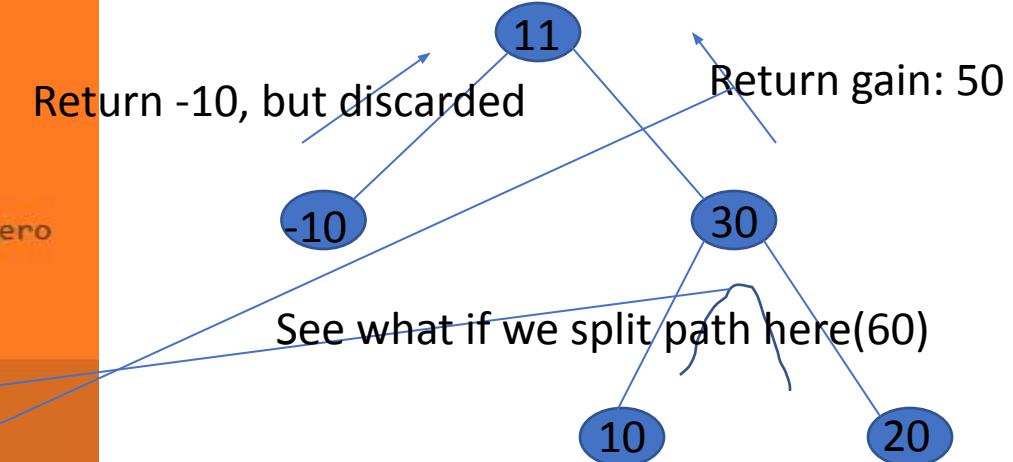
LC string#4.4: 102. Binary Tree PostOrder

Traversal:

```
12 class Solution {
13 public:
14     vector<int> postorderTraversal(TreeNode* root) {
15         vector<int> res;
16         if (root == NULL) return res;
17 #ifdef RECURSIVE
18         vector<int> left = preorderTraversal(root->left);
19         for (int i = 0; i < left.size(); i++)
20             res.push_back(left[i]);
21         vector<int> right = preorderTraversal(root->right);
22         res.push_back(root->val);
23         for (int i = 0; i < right.size(); i++)
24             res.push_back(right[i]);
25 #endif
26
27         stack<TreeNode *> st;
28         st.push(root);
29         while (!st.empty()) {
30             TreeNode* cur = st.top();
31             st.pop();
32             res.insert(res.begin(), cur->val); // this makes all existing one shift to right
33             if (cur->left) st.push(cur->left); // left node first push
34             if (cur->right) st.push(cur->right);
35         }
36
37         return res;
38     }
39 }
```

LC tree#5: 124. Binary Tree Maximum Path Sum(HARD):

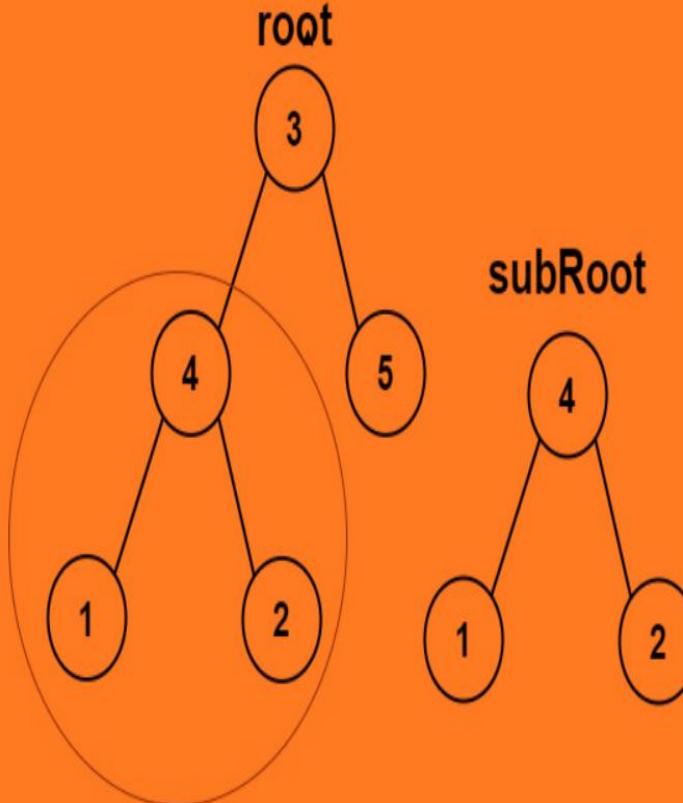
```
12 * class Solution {
13     int res = INT_MIN;
14 public:
15     // return gain of node without split path
16     int max_gain(TreeNode *n) {
17         if (n == NULL) return 0;
18
19         // if gain is negative, we will not take it, so set to zero
20         int left_gain = max(max_gain(n->left), 0);
21         int right_gain = max(max_gain(n->right), 0);
22
23         // check what is value if we split path at this node
24         if (n->val + left_gain + right_gain > res)
25             res = n->val + left_gain + right_gain;
26
27         return n->val + max(left_gain, right_gain);
28     }
29
30     int maxPathSum(TreeNode* root) {
31         if (root == NULL) return 0;
32
33         max_gain(root);
34
35         return res;
36     }
37 };
```



LC tree#6:

572. Subtree of Another Tree:

Example 1:



Input: root = [3,4,5,1,2], subRoot = [4,1,2]

Output: true

- Generate list of nodes starting from root
- Using each node to do same tree checking
- Same tree check using traversal to generate vector<int> to compare.

```
12 class Solution {  
13     public:  
14         vector<int> levelTraver(TreeNode* root) {  
15             vector<int> res;  
16             if (root == NULL) return res;  
17             queue<TreeNode *> Q;  
18             Q.push(root);  
19             TreeNode *dummy = new TreeNode(INT_MAX);  
20             while (!Q.empty()) {  
21                 TreeNode *n = Q.front(); Q.pop();  
22                 res.push_back(n->val);  
23                 if (n->val != INT_MAX) {  
24                     Q.push(n->left ? n->left:dummy);  
25                     Q.push(n->right ? n->right:dummy);  
26                 }  
27             }  
28             return res;  
29         }  
30         bool sameTree(TreeNode *p, TreeNode *q) {  
31             vector<int> p_v = levelTraver(p);  
32             vector<int> q_v = levelTraver(q);  
33             return p_v == q_v;  
34         }  
35         bool isSubtree(TreeNode* root, TreeNode* subRoot) {  
36             vector<TreeNode *> nodes;  
37             if (root == NULL) return false;  
38             queue<TreeNode *> Q;  
39             Q.push(root);  
40             while (!Q.empty()) {  
41                 TreeNode *n = Q.front(); Q.pop();  
42                 nodes.push_back(n);  
43                 if (n->left) Q.push(n->left);  
44                 if (n->right) Q.push(n->right);  
45             }  
46             for (int i = 0; i < nodes.size(); i++) {  
47                 if (nodes[i]->val == subRoot->val) {  
48                     if (sameTree(nodes[i], subRoot))  
49                         return true;  
50                 }  
51             }  
52         }  
53     }  
54 }
```

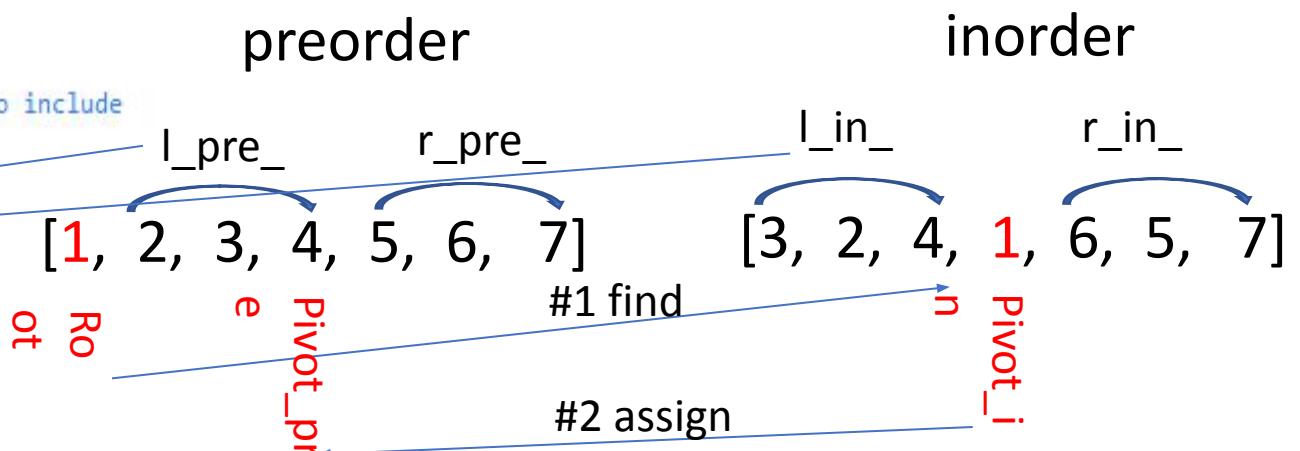
LC tree#7: 105. Construct Binary Tree from Preorder and Inorder Traversals:

```
12+ class Solution {  
13 public:  
14     TreeNode* buildTree(vector<int>& preorder, vector<int>& inorder) {  
15         //first or preorder is for root  
16         if (preorder.size() == 0) return NULL;  
17  
18         TreeNode *root = new TreeNode(preorder[0]); if (preorder.size() == 1) return root;  
19         // find the pivot position of root node inside inorder  
20         // (because of unique value) using root value  
21         int pivot_in = find(inorder.begin(), inorder.end(), preorder[0]) - inorder.begin(); #1  
22         //partition preorder, use pivot_in (inclusive)  
23         int pivot_pre = pivot_in; #2  
24         vector<int> l_pre_, r_pre_;  
25         for (int j = 1; j <= pivot_pre; j++) // exclude root , start with 1  
26             l_pre_.push_back(preorder[j]);  
27         for (int j = pivot_pre+1 ; j < preorder.size(); j++)  
28             r_pre_.push_back(preorder[j]);  
29         //partition inorder, pivot is pivot_in  
30         vector<int> l_in_, r_in_;  
31         for (int j = 0; j < pivot_in; j++) // pivot_in is root, not to include  
32             l_in_.push_back(inorder[j]);  
33         for (int j = pivot_in+1; j < inorder.size(); j++)  
34             r_in_.push_back(inorder[j]);  
35  
36         root->left = buildTree(l_pre_, l_in_);  
37         root->right = buildTree(r_pre_, r_in_);  
38  
39         return root;  
40     }  
41 }
```

Given peorder and inorder, this precisely define a tree
First value from preorder is ROOT, which can be used
To find the partition of LEFT and RIGHT!
If only preorder/inorder given, unless it is a full balanced tree
If postorder is given instead of preorder, find root using last!

Time & Space: O(n)

#1



LC tree#8: 226. serialize and deserialize binary tree:

```
10
11     class Codec {
12
13     // Encodes a tree to a single string following preorder sequence
14     string serialize(TreeNode* root) {
15         if (root == NULL) return "9999,"; // flag as end
16
17         string res = to_string(root->val) + ",";
18         res.append(serialized(root->left)));
19         res.append(serialized(root->right));
20         return res;
21     }
22     int pos = 0; // use to track current processed node
23     TreeNode * buildTree(vector<int>& vals) {
24         if (vals[pos] == 9999) {
25             pos++; return NULL;
26         }
27         TreeNode * root = new TreeNode(vals[pos]); pos++;
28         root->left = buildTree(vals);
29         root->right = buildTree(vals);
30         return root;
31     }
32     // Decodes your encoded data to tree.
33     TreeNode* deserialize(string data) {
34         vector<int> vals;
35         int s = 0, len = 0;
36         for (int i = 0; i < data.size(); i++) {
37             if (data.at(i) != ',')
38                 len++;
39             else {
40                 vals.push_back(stoi(data.substr(s, len)));
41                 s = i+1;
42                 len = 0;
43             }
44         }
45         return buildTree(vals);
46     }
47 };
```

Use simple recursion/dfs to do preorder traversal
Using 9999 as NULL node.
Processing serialized string into `vector<int>` and use
Recursion to deserialize tree node.
Use variable “**pos**” to track position.

LC tree#9: 98. Validate Binary Search Tree:

```
12 * class Solution {
13 public:
14     bool helper (TreeNode *root, int64_t low, int64_t high) {
15         if (root->val > low && root->val < high) {
16             if (root->left)
17                 if (!helper(root->left, low, root->val)) return false;
18             if (root->right)
19                 if (!helper(root->right, root->val, high)) return false;
20         } else {
21             return false;
22         }
23         return true;
24     }
25     bool isValidBST(TreeNode* root) {
26         if (root == NULL) return false;
27         int64_t LOW_INF = ((int64_t)INT_MIN)-1;
28         int64_t HIGH_INF = ((int64_t)INT_MAX) + 1;
29         return helper(root, LOW_INF, HIGH_INF);
30     }
31 }
```

- Use `int64_t LOW_INF = (int64_t)INT_MIN-1`
- Use `int64_t HIGH_INF = (int64_t)INT_MAX+1`
- To check left node, need pass down current node value as HIGH
- To check right node, need pass down current node value as LOW
- If there is any false return, return back all the way.

LC tree#10: 230. Kth Smallest Element in a BST:

```
12 class Solution {  
13     public:  
14         void inorderT(TreeNode * root, vector<int>& list) {  
15             if (!root) return;  
16             if (root->left) inorderT(root->left, list);  
17             list.push_back(root->val);  
18             if (root->right) inorderT(root->right, list);  
19         }  
20         int kthSmallest(TreeNode* root, int k) {  
21             vector<int> list;  
22             inorderT(root, list);  
23             return list[k-1];  
24         }  
25     };
```

BST inorder traversal will give list of Value in ascending order.

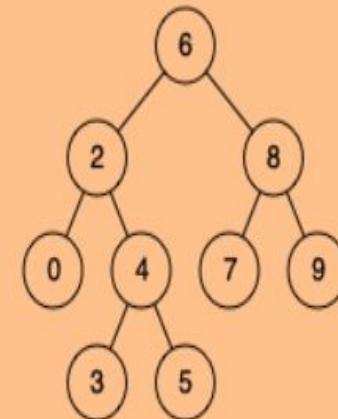
```
int height_dfs(Node* node, int &res) {  
    if (node == NULL) return -1;  
    if (node->children.size() == 0) return 0;  
  
    vector<int> heights;  
    for (auto n: node->children) {  
        heights.push_back(height_dfs(n, res));  
    }  
    sort(heights.begin(), heights.end(), greater<int>());  
    if (heights.size() >= 2) {  
        res = max(res, heights[0] + heights[1] + 2); // add two: one for each side  
    }  
    return heights[0]+1;  
}
```

LC tree#11: 235. Lowest Common Ancestor of a Binary Search Tree:

- Find p in BST into vector<int>
- Find q in BST into unordered_map<int, TreeNode*>

```
11 class Solution {
12 public:
13     void findNodeList(TreeNode * root, TreeNode *p, vector<int> &list) {
14         TreeNode *cur = root;
15         while (cur && cur->val != p->val) {
16             list.insert(list.begin(), cur->val);
17             (cur->val > p->val)? (cur = cur->left):(cur = cur->right);
18         }
19         list.insert(list.begin(), cur->val);
20     }
21     void findNodeMap(TreeNode * root, TreeNode *p, unordered_map<int, TreeNode*> &map) {
22         TreeNode *cur = root;
23         while (cur && cur->val != p->val) {
24             map[cur->val] = cur;
25             (cur->val > p->val)? (cur = cur->left):(cur = cur->right);
26         }
27         map[cur->val] = cur;
28     }
29     TreeNode* lowestCommonAncestor(TreeNode* root, TreeNode* p, TreeNode* q) {
30         vector<int> p_list;
31         unordered_map<int, TreeNode*> q_map;
32         findNodeList(root, p, p_list);
33         findNodeMap(root, q, q_map);
34         for (int i = 0; i < p_list.size(); i++){
35             if (q_map.find(p_list[i]) != q_map.end())
36                 return q_map[p_list[i]];
37         }
38         return NULL;
39     }
40 };
```

Example 1:



Input: root = [6,2,8,0,4,7,9,null,null,3,5], p = 2, q = 8

Output: 6

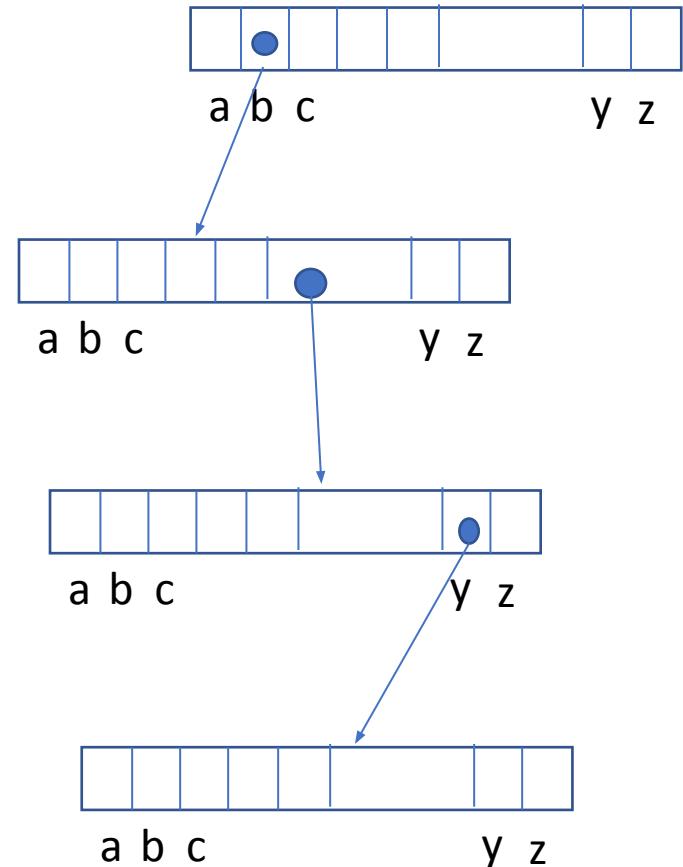
Explanation: The LCA of nodes 2 and 8 is 6.

LC tree#12: 208. Implement Trie (Prefix Tree)

```
1  class Trie {
2     Trie* children[26];
3     bool endOfWord;
4
5 public:
6     /** Initialize your data structure here. */
7     Trie() {
8         for (int i = 0; i < 26; i++) children[i] = NULL;
9         endOfWord = false;
10    }
11    /** Inserts a word into the trie. */
12    void insert(string word) {
13        Trie *p = this;
14        for (int i = 0; i < word.length(); i++) {
15            char c = word.at(i); // each character is used as key into children
16            if (p->children[c-'a'] == NULL) {
17                Trie *t = new Trie();
18                p->children[c-'a'] = t;
19            }
20            p = p->children[c-'a'];
21        }
22        p->endOfWord = true;
23    }
24    /** Returns if the word is in the trie. */
25    bool search(string word) {
26        Trie *p= this;
27        for (int i = 0; i < word.length(); i++) {
28            char c = word.at(i);
29            p = p->children[c-'a'];
30            if (!p) return false;
31        }
32        return p->endOfWord;
33    }
34    /** Returns if there is any word in the trie that starts with the given prefix. */
35    bool startsWith(string prefix) {
36        Trie *p= this;
37        for (int i = 0; i < prefix.length(); i++) {
38            char c = prefix.at(i);
39            p = p->children[c-'a'];
40            if (!p) return false;
41        }
42        return (p != NULL);
43    }
44};
```

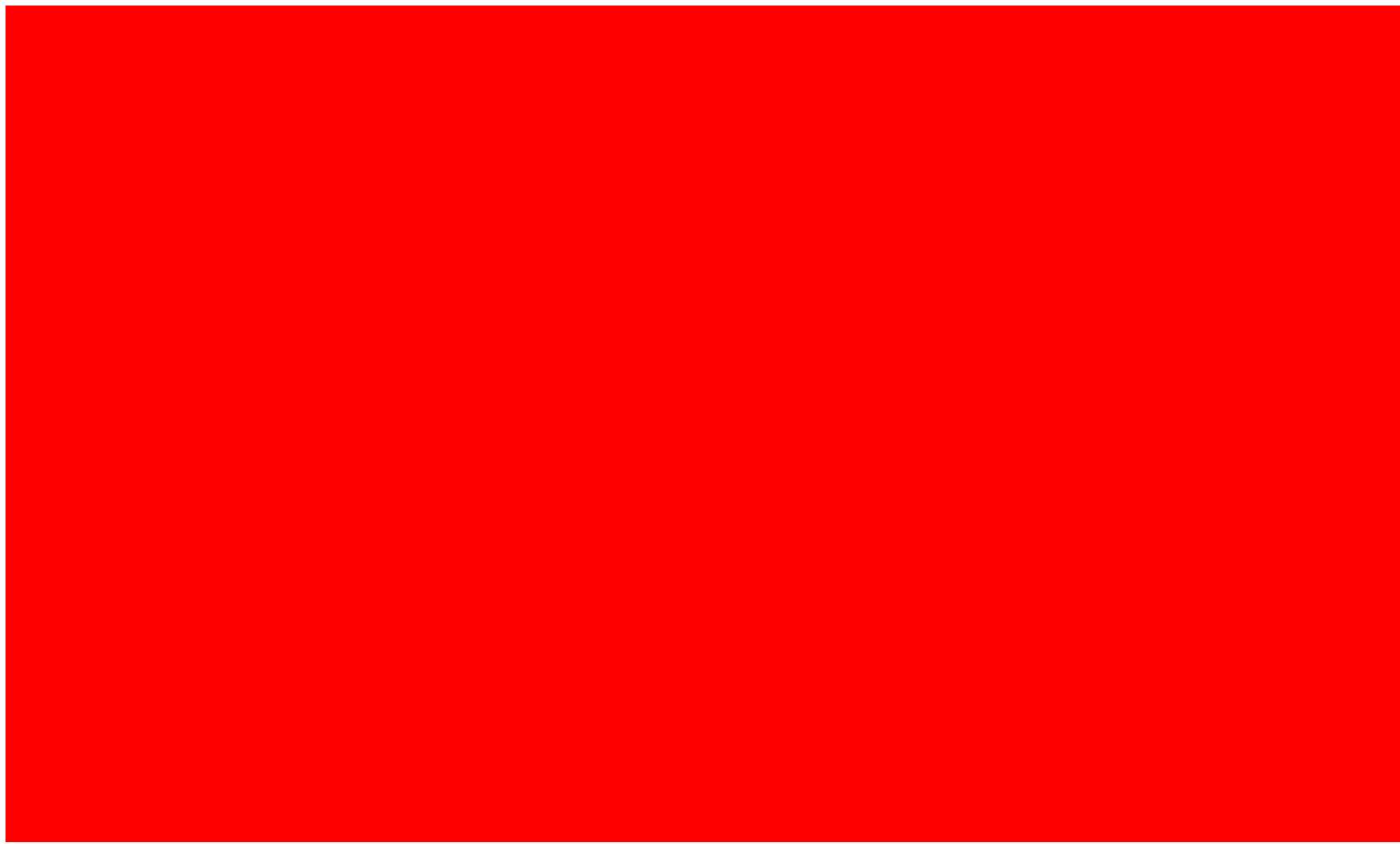
Insert("boy")

ROOT



LC tree#13: 211. Design Add and Search Words Data Structure:
(essential trie)

LC tree#14: 212. Word Search II: (HARD)



LC Array#1: 1. Two Sum (easy)

Using unordered_map<int, int> to record
Index and value for search: O(n)

```
1  class Solution {
2  public:
3      vector<int> twoSum(vector<int>& nums, int target) {
4          unordered_map<int, int> umap;
5          unordered_map<int, int>::iterator itr;
6          vector<int> res;
7
8          for (int i = 0; i < nums.size(); i++) {
9              int remainder = target - nums[i];
10             itr = umap.find(remainder);
11             if (itr != umap.end()) {
12                 res.push_back(i);
13                 res.push_back(itr->second);
14                 break;
15             } else {
16                 umap[nums[i]] = i;
17             }
18         }
19         return res;
20     }
21 };
```

LC Array#2: 121. Best Time to Buy and Sell Stock (I), only one transaction

122: ((II) allow to buy and sell same day, no limit of transactions

```
1 class Solution {  
2     public:  
3         int maxProfit(vector<int>& prices) {  
4             int cost = prices[0];  
5             int profit = 0;  
6             for (int i = 1; i < prices.size(); i++) {  
7                 if (prices[i] >= cost) {  
8                     profit = max(profit, prices[i] - cost);  
9                 } else {  
10                     cost = prices[i]; // starting over if we see lower co  
11                 }  
12             }  
13             return profit;  
14         }  
15     };
```

```
1 class Solution {  
2     public:  
3         int maxProfit(vector<int>& prices) {  
4             int profit = 0;  
5             for (int i = 1; i < prices.size(); i++) {  
6                 if (prices[i] > prices[i-1]) { // for any day, if we seen today price is higher  
7                     profit += prices[i] - prices[i-1]; // we always say we bought it yesterday  
8                 } // if tomorrow is even higher, we said we bought it back yesterday  
9             }  
10            return profit;  
11        }  
12    };
```

- a

LC Array#3: 217. Contains Duplicate (easy) 219: Contains Duplicate II

```
1 class Solution {
2 public:
3     bool containsDuplicate(vector<int>& nums) {
4         unordered_set<int> seen;
5         for (int i = 0; i < nums.size(); i++) {
6             if (seen.find(nums[i]) == seen.end()) {
7                 seen.insert(nums[i]);
8             } else {
9                 return true;
10            }
11        }
12        return false;
13    }
14};
```

Example 1:

```
Input: nums = [1,2,3,1]
Output: true
```

- Scan thru list, use hashSet to record if not seen

Example 2:

```
Input: nums = [1,2,3,4]
Output: false
```

```
1 class Solution {
2 public:
3     bool containsNearbyDuplicate(vector<int>& nums, int k) {
4         unordered_map<int, int> map;
5         for (int i = 0; i < nums.size(); i++) {
6             if (map.find(nums[i]) == map.end()) {
7                 map[nums[i]] = i;
8             } else {
9                 int j = map[nums[i]];
10                if (i - j <= k)
11                    return true;
12                else
13                    map[nums[i]] = i;
14            }
15        }
16        return false;
17    }
18};
```

LC Array#4: 238. Product of Array Except Self

```
1 class Solution {
2 public:
3     vector<int> productExceptSelf(vector<int>& nums) {
4         vector<int> left(nums.size()), right(nums.size());
5         int tmp = 1;
6         left[0] = tmp;
7         for (int i=1; i < nums.size(); i++){
8             tmp *= nums[i-1];
9             left[i] = tmp;
10        }
11        tmp = 1;
12        right[nums.size()-1] = tmp;
13        for (int i = nums.size()-2; i >= 0; i--){
14            tmp *= nums[i+1];
15            right[i] = tmp;
16        }
17
18        vector<int> res(nums.size());
19        for (int i = 0; i < nums.size(); i++) {
20            res[i] = left[i]*right[i];
21        }
22        return res;
23    }
24};
```

- One pass starting from beginning to make LEFT product
- Another pass starting from end to make RIGHT product

LC Array#5: 53. Maximum Subarray

```
1 * class Solution {
2 public:
3     int maxSubArray(vector<int>& nums) {
4         int maxSum = INT_MIN;
5         int curSum = 0;
6         for (int i = 0; i < nums.size(); i++) {
7             curSum += nums[i];
8             maxSum = max(maxSum, curSum);
9             if (curSum < 0) curSum = 0;
10        }
11
12        return maxSum;
13    }
14}
```

- Use two variables: curSum, maxSum
- Reset to curSum back to zero if it is negative

LC Array#6: 152. Maximum Product Subarray

```
1 * class Solution {
2 public:
3     int maxProduct(vector<int>& nums) {
4         int maxRes = INT_MIN;
5
6         int product = 1;
7         for (int i = 0; i < nums.size(); i++) { //left to right
8             product *= nums[i];
9             maxRes = max(product, maxRes);
10            if (product == 0) product = 1;
11        }
12        product = 1;
13        for (int i = nums.size()-1; i >= 0; i--) { // right to left|
14            product *= nums[i];
15            maxRes = max(product, maxRes);
16            if (product == 0) product = 1;
17        }
18
19        return maxRes;
20    }
21};
```

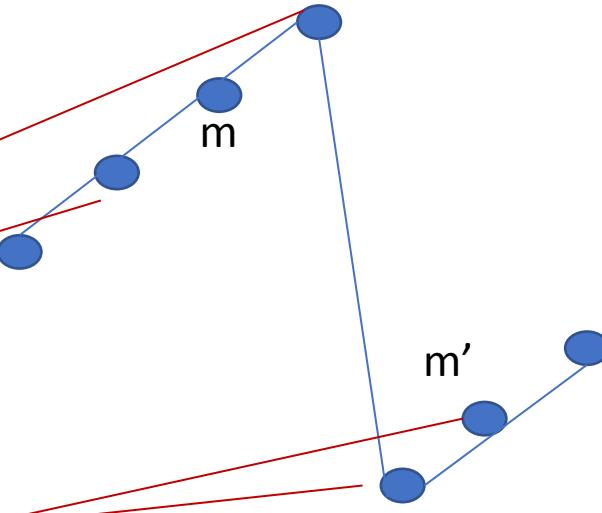
- Use two variables: maxRes, product
- Do product from left to right, track maxRes, if product becomes 0, reset to 1
- Do product from right to left, track maxRes, if product becomes 0, reset to 1

Binary search: three forms

int

LC Array#7: 153. Find Minimum in Rotated Sorted Array in $O(\log n)$ time.

```
1 class Solution {
2     public:
3         int findMin(vector<int>& nums) {
4             if (nums.size() == 1) return nums[0];
5             if (nums[0] < nums[nums.size()-1]) return nums[0];
6
7             int l = 0, r= nums.size() - 1;
8             while (l <= r) {
9                 if (r == 1) return nums[l]; // special handling
10                if (r-1 ==1) return (min(nums[l], nums[r])); // special handling
11
12                int m = l + (r-1)/2;
13                if (nums[m] > nums[0]) {
14                    if (m+1 < nums.size() && nums[m] > nums[m+1]) {
15                        return nums[m+1];
16                    } else {
17                        l = m + 1;
18                    }
19                } else {
20                    if (m-1 >= 0 && nums[m-1] > nums[m]) {
21                        return nums[m];
22                    } else {
23                        r = m -1;
24                    }
25                }
26            }
27            return INT_MIN;
28        }
29    };
}
```



- Special handling when $r=l$ and $r-l == 1$ inside loop
- Special handling when $\text{nums}[0]$ is minimum

LC Array#7.1: 162. Find Peak Element in $O(\log n)$ time.

```
1 class Solution {  
2     public:  
3         int findPeakElement(vector<int>& nums) {  
4             int l = 0, r = nums.size() - 1;  
5             if (r == 0) return 0;  
6             while (l < r) {  
7                 if (r-1 == 1) return nums[1] > nums[r]? 1:r;  
8                 int m = (l+r)/2;  
9                 if (nums[m] < nums[m+1]) // climbing  
10                     l = m;  
11                 else  
12                     r = m;  
13             }  
14             return l;  
15         }  
16     };
```

Input: nums = [1,2,1,3,5,6,4]

Output: 5

Explanation: Your function can return either index number 1 where the peak element is 2, or index number 5 where the peak element is 6.

A peak element is an element that is strictly greater than its neighbors.

Given an integer array `nums`, find a peak element, and return its index. If the array contains multiple peaks, return the index to **any of the peaks**.

You may imagine that `nums[-1] = nums[n] = -∞`.

You must write an algorithm that runs in $O(\log n)$ time.

LC Array#8: 33. Search in Rotated Sorted Array using O(log N)

```
1 class Solution {
2     public:
3         int bst(vector<int>& nums, int l, int r, int target) {
4             int res = -1;
5             while (l <= r) {
6                 int m = (l+r)/2;
7                 if (nums[m] == target) {
8                     res = m; break;
9                 } else {
10                     (nums[m] < target) ? (l=m+1):(r=m-1);
11                 }
12             }
13             return res;
14         }
15         int search(vector<int>& nums, int target) {
16             int pivot, l = 0 , r = nums.size()-1;
17             if (nums.size() == 1)
18                 return nums[0] == target? 0:-1;
19             if (nums[l] < nums[r]) return bst(nums, l, r, target);
20             while (l<=r) {
21                 if (l == r) {
22                     pivot = l; break;
23                 }
24                 if (r-l == 1) {
25                     pivot = nums[l] < nums[r]? l:r;
26                     break;
27                 }
28                 int m = (l+r)/2;
29                 if (m>0 && (m+1) < (nums.size()) && nums[m+1] > nums[m] && nums[m-1] > nums[m]) {
30                     pivot = m;
31                     break;
32                 }
33                 if (m>0 && (m+1) < (nums.size()) && nums[m-1] < nums[m] && nums[m] > nums[m+1]) {
34                     pivot = m+1;
35                     break;
36                 }
37                 if (nums[m] > nums[0])
38                     l = m+1;
39                 else
40                     r = m -1;
41             }
42             if (target >= nums[0])
43                 return bst(nums, 0, pivot-1, target);
44             else
45                 return bst(nums, pivot, nums.size()-1, target);
46         };
47     }
```

- If the array size is 1, return result
- If array is fully sorted, do binary search
- Otherwise find pivot, then decide which segment do perform binary search.

LC Array#9: 15. 3Sum.

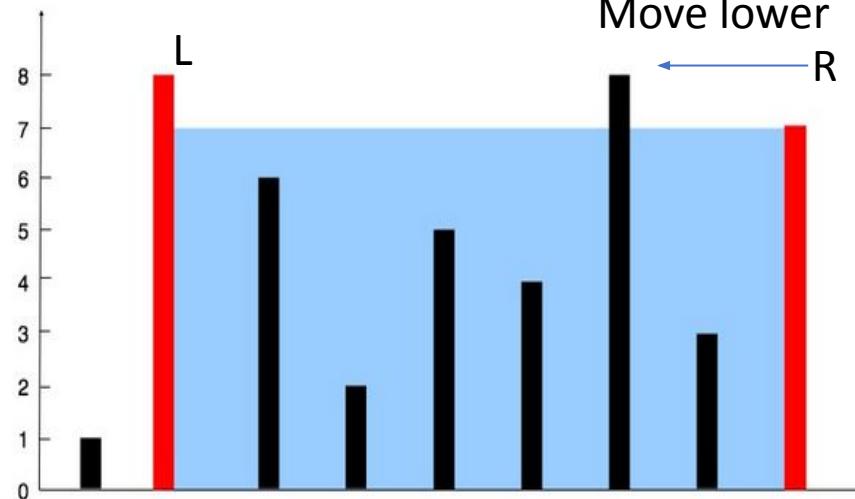
```
1 class Solution {
2 public:
3     vector<vector<int>> threeSum(vector<int>& nums) {
4         vector<vector<int>> res;
5         if (nums.size() < 3) return res;
6         sort(nums.begin(), nums.end());
7         for (int i = 0; i < nums.size()-2; i++) {
8             if (nums[i] > 0) continue;
9             if (i > 0 && (nums[i] == nums[i-1])) continue; // avoid dup
10            int target = -nums[i];
11            int l = i+1; int r = nums.size()-1;
12            while (l < r) {
13                if (nums[l] + nums[r] == target) {
14                    vector<int> ans;
15                    ans.push_back(nums[i]);
16                    ans.push_back(nums[l]);
17                    ans.push_back(nums[r]);
18                    res.push_back(ans);
19                    // don't break, maybe more answer
20                }
21                // advance l or r
22                if (nums[l] + nums[r] <= target) {
23                    int ll = l;
24                    while ((l < nums.size() -1) && (nums[ll] == nums[l])) l++;
25                } else {
26                    int rr = r;
27                    while ((r > 0) && (nums[rr] == nums[r])) r--;
28                }
29            }
30        }
31        return res;
32    }
33};
```

- Sort array
- Go thru for i, target is $-\text{nums}[i]$
- Using two pointers between $i+1$ to $\text{size}()-1$ to find other two numbers
- $O(n^*n)$

LC Array#10: 11. Container With Most Water.

```
1 class Solution {
2 public:
3     int maxArea(vector<int>& height) {
4         int res = INT_MIN;
5
6 #ifdef BRUTEFORCE
7         for (int i = 0; i < height.size() - 1; i++) {
8             for (int j = i+1; j < height.size(); j++) {
9                 int area = (j-i)*min(height[i], height[j]);
10                res = max(res, area);
11            }
12        }
13 #endif
14 // two pointers
15     int l = 0, r = height.size() - 1;
16     while (l < r) {
17         int a = (r-l)*min(height[l], height[r]);
18         res = max(res, a);
19         if (height[l] < height[r]) {
20             l++;
21         } else {
22             r--;
23         }
24     }
25     return res;
26 }
27 }
```

Example 1:



Input: height = [1,8,6,2,5,4,8,3,7]

Output: 49

Explanation: The above vertical lines are represented by array [1,8,6,2,5,4,8,3,7]. In this case, the max area of water (blue section) the container can contain is 49.

LC#528 PrefixSum

528. Random Pick with Weight

Medium 1494 3131 Add to List Share

You are given an array of positive integers `w` where `w[i]` describes the weight of `ith` index (0-indexed).

We need to call the function `pickIndex()` which **randomly** returns an integer in the range `[0, w.length - 1]`. `pickIndex()` should return the integer proportional to its weight in the `w` array. For example, for `w = [1, 3]`, the probability of picking the index `0` is $1 / (1 + 3) = 0.25$ (i.e 25%) while the probability of picking the index `1` is $3 / (1 + 3) = 0.75$ (i.e 75%).

More formally, the probability of picking index `i` is `w[i] / sum(w)`.

```
1 class Solution {
2     vector<int> prefixSum;
3     int N;
4 public:
5     Solution(vector<int>& w) {
6         N = w.size();
7         prefixSum.push_back(w[0]);
8         for (int i = 1; i < N; i++) {
9             prefixSum.push_back(prefixSum.back() + w[i]);
10        }
11    }
12
13    int pickIndex() {
14        float r = ((float)rand()) / RAND_MAX; // get a number between 0-1
15        int target = (int) (r * prefixSum[N-1]);
16        for (int i = 0; i < N; i++) {
17            if (target < prefixSum[i]) return i;
18        }
19        return N-1;
20    }
21 }
22 }
```

Bit Manipulation

10101010001101001101010010010001010
0100010101001010100101010010100101010

Types of Bitwise Operators

Operator	Name	Example	Result
&	Bitwise AND	6 & 3	2
	Bitwise OR	10 10	10
^	Bitwise XOR	2^2	0
~	Bitwise 1's complement	~9	-10
<<	Left-Shift	10<<2	40
>>	Right-Shift	10>>2	2

LC BITS#1: 371. Sum of Two Integers.

```
1 class Solution {  
2     public:  
3         int getSum(int a, int b) {  
4             long mask = 0xFFFFFFFF;  
5             while (b) {  
6                 int sum = a ^ b;  
7                 int carry = ((a & b)&mask) << 1; // avoid negative# shift error  
8                 a = sum;  
9                 b = carry;  
10            }  
11            return a;  
12        }  
13    };  
14  
15
```

- $(A \& B) << 1 \Rightarrow \text{carry}$, $A \wedge B \Rightarrow \text{answer}$
- Use **long** mask=0xFFFF,FFFF

LC BITS#2: 268. Missing Number.

```
1 class Solution {  
2     public:  
3         int missingNumber(vector<int>& nums) {  
4             int res = nums.size();  
5             for (int i = 0; i < nums.size(); i++) {  
6                 res ^= i;  
7                 res ^= nums[i];  
8             }  
9             return res;  
10        }  
11    };
```

Given an array `nums` containing n distinct numbers in the range $[0, n]$, return *the only number in the range that is missing from the array*.

Follow up: Could you implement a solution using only $O(1)$ extra space complexity and $O(n)$ runtime complexity?

Example 1:

Input: `nums = [3,0,1]`

Output: 2

Explanation: $n = 3$ since there are 3 numbers, so all numbers are in the range $[0,3]$. 2 is the missing number in the range since it does not appear in `nums`.

- $(A \& B) << 1 \Rightarrow \text{carry}, A^B \Rightarrow \text{answer}$

LC BITS#3: 190. Reverse Bits.

```
1 class Solution {
2 public:
3     uint32_t reverseBits(uint32_t n) {
4         unsigned int l_mask = 1 << 31, r_mask = 1;
5         for (int i = 0; i < 16; i++) {
6             int r_ = (n & r_mask);
7             int l_ = (n & l_mask);
8             n = (n & ~l_mask) | (r_?l_mask:0);
9             n = (n & ~r_mask) | (l_?r_mask:0);
10            l_mask = l_mask >>1 ;
11            r_mask = r_mask << 1;
12        }
13        return n;
14    }
15};
```

• a

LC BITS#4: counting 1 bit of an integer.

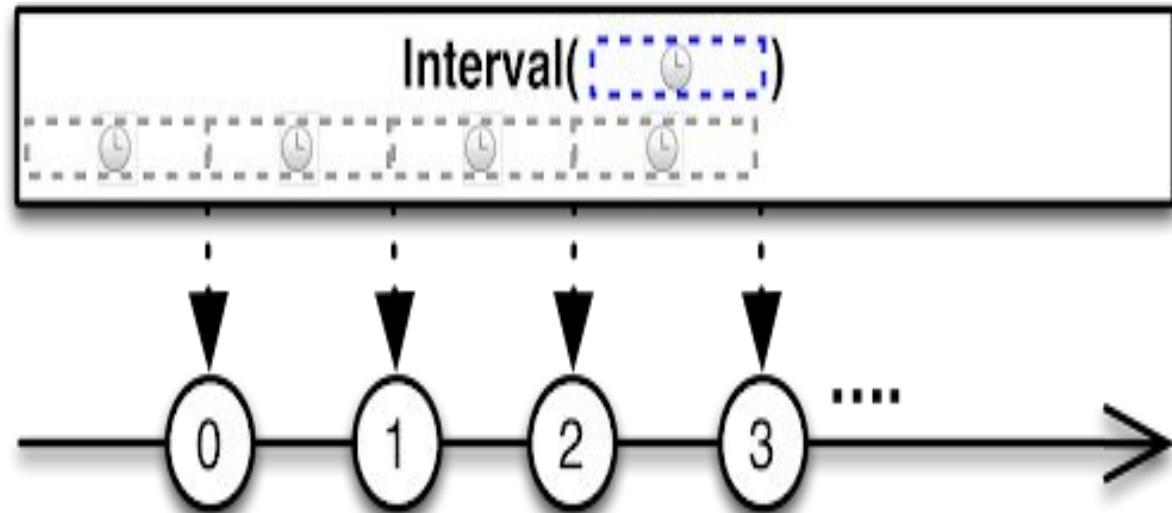
```
while (n) { counter++; n &=(n-1);}  
__builtin_popcount(n);
```

- `__builtin_popcount()`: count number of 1 bit, leverage “**popcnt**” instruction
- `__builtin_clz()`: count leading zero, “**lea**” instruction
- `__builtin_ctz()`: count trailing zero, “**tzcnt**” instruction
- `__builtin_parity()`: parity check

LC BITS#5: check if a number is power of 4.

```
bool isPowerOfFour(int n) {  
    return !(n&(n-1)) && (n&0x55555555);  
    //check the 1-bit location;  
}
```

- Check it has only one bit and bit location is 0x0101, 0101, 0101, 0101, 0101, 0101, 0101
- If it is power of 8, which is $(2 \times 2 \times 2)N$, bit location:
0x0100, 1001, 0010, 0100, 1001, 0010, 0100, 1001
0x49249249



LC INTERVAL#1: 57. Insert Interval.

```
1 class Solution {
2 public:
3     vector<vector<int>> insert(vector<vector<int>> & intervals, vector<int>& newInterval) {
4         //insert newInterval into intervals to keep ascending order
5         bool inserted = false;
6         for (int i = 0; i < intervals.size(); i++) {
7             if (newInterval[0] <= intervals[i][0]) {
8                 intervals.insert(intervals.begin() + i, newInterval);
9                 inserted = true;
10                break;
11            }
12        }
13        if (!inserted) intervals.push_back(newInterval);
14    }
15
16    stack<vector<int>> st;
17    st.push(intervals[0]);
18    for (int i = 1; i < intervals.size(); i++) {
19        vector<int> top = st.top();
20        if (intervals[i][0] > top[1]) { // no-overlap,
21            st.push(intervals[i]);
22        } else {
23            if (intervals[i][1] > top[1]) { // bigger than current interval
24                top[1] = intervals[i][1]; // extend current interval
25                st.pop();
26                st.push(top);
27            }
28        }
29    }
30    vector<vector<int>> res;
31    while (!st.empty()) {
32        res.insert(res.begin(), st.top());
33        st.pop();
34    }
35}
36};
```

Example 1:

Input: intervals = [[1,3],[6,9]], newInterval = [2,5]
Output: [[1,5],[6,9]]

- Given interval list is in ascending order
- Insert newInterval into given list in ascending order
- Merge all intervals in the list using **STACK**

LC INTERVAL#2: 56. Merge Intervals

```
1 class Solution {
2 public:
3     static bool cmp(vector<int> v1, vector<int> v2) {
4         return v1[0] < v2[0];
5     }
6     vector<vector<int>> merge(vector<vector<int>> intervals) {
7         if (intervals.size() == 0) return res;
8         sort(intervals.begin(), intervals.end(), cmp);
9
10    vector<vector<int>> res;
11    stack<vector<int>> st;
12    st.push(intervals[0]);
13    for (int i = 1; i < intervals.size(); i++) {
14        vector<int> top = st.top();
15        int c_s = top[0];
16        int c_e = top[1];
17        int n_s = intervals[i][0];
18        int n_e = intervals[i][1];
19        if (n_s <= c_e) {
20            // merge
21            top[1] = n_e > c_e ? n_e: c_e;
22            st.pop();
23            st.push(top);
24        } else {
25            // no overlap, simply push
26            st.push(intervals[i]);
27        }
28    }
29
30    while (!st.empty()) {
31        res.insert(res.begin(), st.top());
32        st.pop();
33    }
34    return res;
35}
36};
```

Example 1:

Input: intervals = [[1,3],[2,6],[8,10],[15,18]]

Output: [[1,6],[8,10],[15,18]]

Explanation: Since intervals [1,3] and [2,6] overlaps, merge them into [1,6].

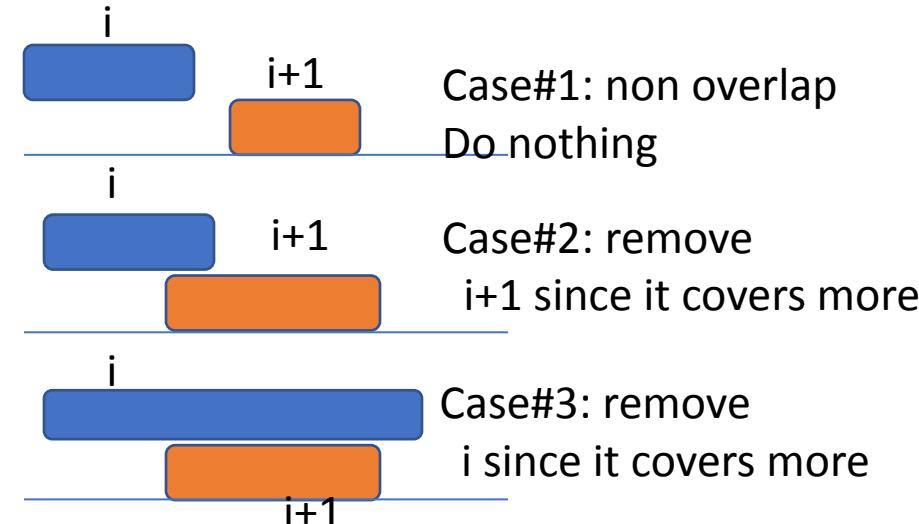
- First sort intervals by using start value in ascending order
- Use stack to merge any overlapped intervals

LC INTERVAL#3: 435. Non-overlapping Intervals.

Given an array of intervals `intervals` where `intervals[i] = [starti, endi]`, return the minimum number of intervals you need to remove to make the rest of the intervals non-overlapping.

```
1 class Solution {
2     public:
3         static bool cmp(vector<int> v1, vector<int> v2) {
4             return v1[0] < v2[0];
5         }
6         // first sort the intervals based on starting value in ascending order
7         // if interval[i] overlaps with interval[i+1]
8         //   if interval[i+1] end is greater than interval[i], remove i+1
9         //   else remove interval[i], which is longer
10        int eraseOverlapIntervals(vector<vector<int>>& intervals) {
11            if (intervals.size() < 2) return 0;
12
13            sort(intervals.begin(), intervals.end(), cmp);
14            vector<int> prev = intervals[0];
15            int removals = 0;
16            for (int i = 1; i < intervals.size(); i++) {
17                vector<int> cur = intervals[i];
18                if (cur[0] >= prev[1]) {
19                    prev = cur;
20                    continue; // non overlap
21                } else {
22                    if (cur[1] > prev[1]) {
23                        // remove cur
24                        removals++;
25                        // prev pointer stays
26                    } else {
27                        prev = cur;
28                        removals++;
29                    }
30                }
31            }
32            return removals;
33        }
34    };
```

- Sort all interval lists in ascending order
- Go thru interval list, there are three cases below:



LC INTERVAL#4: 252. Meeting Rooms.

Given an array of meeting time intervals where $\text{intervals}[i] = [\text{start}_i, \text{end}_i]$, determine if a person could attend all meetings.

```
1 class Solution {
2     public:
3         static bool cmp(vector<int> v1, vector<int> v2) {
4             return v1[0]<v2[0];
5         }
6
7         bool canAttendMeetings(vector<vector<int>>& intervals) {
8             if (intervals.size() < 2) return true;
9
10            sort(intervals.begin(), intervals.end(), cmp);
11
12            vector<int> prev = intervals[0];
13            for (int i = 1; i < intervals.size(); i++) {
14                vector<int> cur = intervals[i];
15                if (prev[1] > cur[0])
16                    return false;
17                else {
18                    prev = cur;
19                }
20            }
21
22            return true;
23        }
24    };
```

Example 1:

Input: intervals = [[0,30],[5,10],[15,20]]
Output: false

- Sort interval by starting time in ascending order
- Check if there is any overlap

LC INTERVAL#5: 253: Meeting Rooms II.

Given an array of meeting time intervals intervals where intervals[i] = [starti, endi], return the minimum number of conference rooms required.

```
1 class Solution {
2     public:
3         static bool cmp(int i1, int i2) {
4             if (abs(i1) == abs(i2)) {
5                 return i1 < i2;
6             }
7             return abs(i1) < abs(i2);
8         }
9
10        int minMeetingRooms(vector<vector<int>>& intervals) {
11            vector<int> times;
12            for (int i = 0; i < intervals.size(); i++) {
13                times.push_back(intervals[i][0]);
14                times.push_back(-intervals[i][1]);
15            }
16            sort(times.begin(), times.end(), cmp);
17            int res = 0;
18            int rooms = 0;
19            for (int i = 0; i < times.size(); i++) {
20                if (times[i] >= 0) {
21                    rooms++;
22                    res = rooms > res ? rooms : res;
23                } else {
24                    rooms--;
25                }
26            }
27        }
28    };
29 }
```

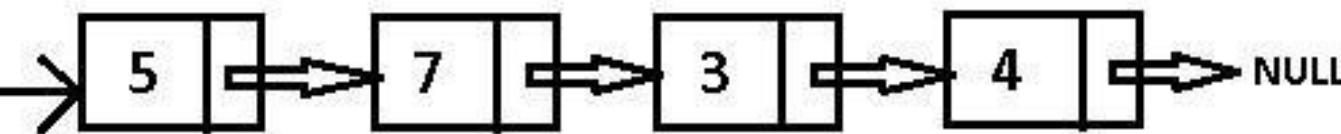
Example 1:

Input: intervals = [[0,30],[5,10],[15,20]]
Output: 2

- Merge all starting time and ending time into one single vector, but ending time as negative value
- Sort vector using absolute value, if it is equal, negative first
- Go thru vector list, see positive time, increase 1,
- See negative time, decrease one.
- Record the max rooms

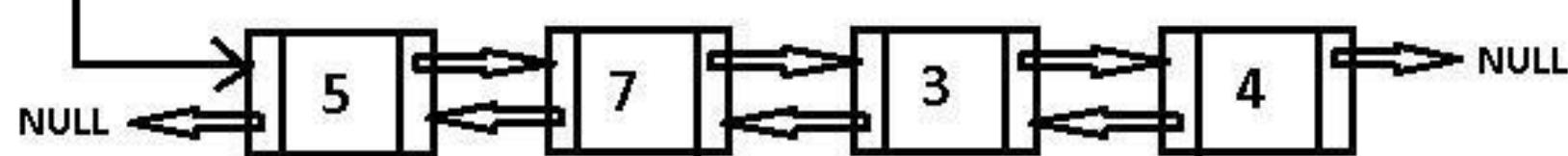
HEAD

Single Linked List



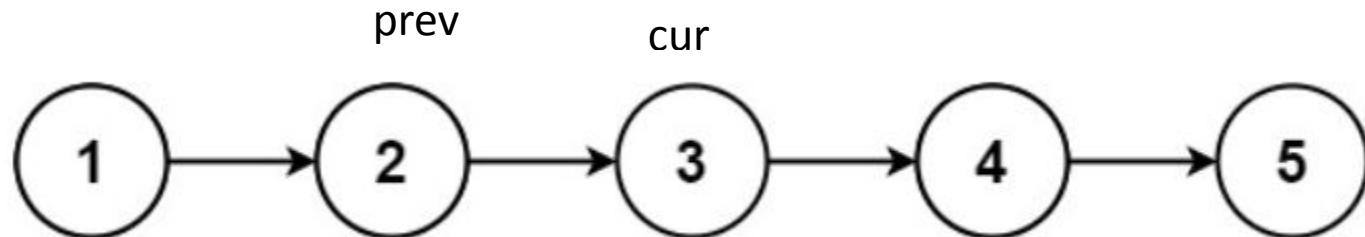
HEAD

Double Linked List



LC LinkedList#1: 206. Reverse Linked List

```
1  /**
2  * Definition for singly-linked list.
3  * struct ListNode {
4  *     int val;
5  *     ListNode *next;
6  *     ListNode() : val(0), next(nullptr) {}
7  *     ListNode(int x) : val(x), next(nullptr) {}
8  *     ListNode(int x, ListNode *next) : val(x), next(next) {}
9  * };
10 */
11 class Solution {
12 public:
13     ListNode* reverseList(ListNode* head) {
14         ListNode *prev = NULL, *cur = head;
15         while (cur) {
16             ListNode *tmp = cur->next;
17             cur->next = prev;
18             prev = cur;
19             cur = tmp;
20         }
21         return prev;
22     }
23 }
```



LC LinkedList#2: 141. Linked List Cycle

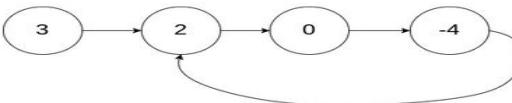
```
9  class Solution {
10 public:
11     bool hasCycle(ListNode *head) {
12         ListNode *s1 = head, *s2;
13         if (!s1) return false;
14         if (s1->next) s2 = s1->next;
15         bool cycle = false;
16         while (s1 && s2) {
17             if (s1 == s2) {
18                 cycle = true;
19                 break;
20             }
21             s1 = s1->next;
22             s2 = s2->next? s2->next->next:NULL;
23         }
24     return cycle;
25 }
26 };
```

Given `head`, the head of a linked list, determine if the linked list has a cycle in it.

There is a cycle in a linked list if there is some node in the list that can be reached again by continuously following the `next` pointer. Internally, `pos` is used to denote the index of the node that tail's `next` pointer is connected to. **Note that pos is not passed as a parameter.**

Return `true` if there is a cycle in the linked list. Otherwise, return `false`.

Example 1:



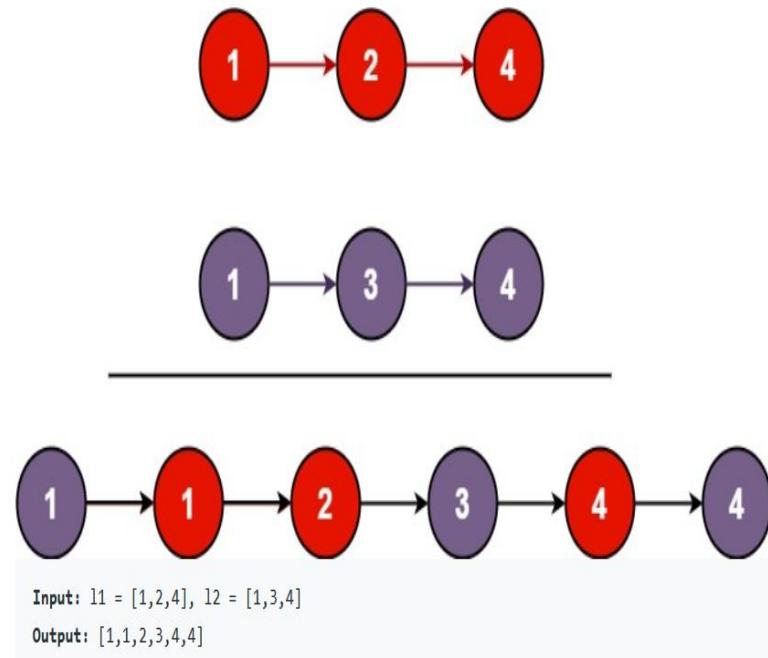
- Use two pointers: `s1` move with one step, `s2` move two steps
- If there is loop, `s1` and `s2` will never NULL and will be `s1==s2` eventually.
- Or use **hashSet “seen”**

```
9  class Solution {
10 public:
11     ListNode *detectCycle(ListNode *head) {
12         if (head == NULL) return NULL;
13         unordered_set<ListNode *> seen;
14         ListNode *s1 = head;
15         while (s1) {
16             if (seen.find(s1) != seen.end()) return s1;
17             seen.insert(s1);
18             s1 = s1->next;
19         }
20     return NULL;
21 }
22 };
```

LC LinkedList#3: 21. Merge Two Sorted Lists.

```
11 class Solution {
12 public:
13     ListNode* mergeTwoLists(ListNode* l1, ListNode* l2) {
14         ListNode *head = NULL, *tmp, *cur;
15         if (!l1) return l2;
16         if (!l2) return l1;
17
18         while (l1 && l2) {
19             if (l1->val <= l2->val) {
20                 tmp = l1;
21                 l1 = l1->next;
22             } else {
23                 tmp = l2;
24                 l2 = l2->next;
25             }
26             if (!head) {
27                 head = tmp;
28                 cur = tmp;
29             } else {
30                 cur->next = tmp;
31                 cur = tmp;
32             }
33         }
34         if (l1) {
35             cur->next = l1;
36         }
37         if (l2) {
38             cur->next = l2;
39         }
40
41         return head;
42     }
43 }
```

Example 1:



- Use “cur” to hold current latest merged node
- Use “head” to hold return value
- Compare l1 and l2, advance the smaller value pointer
- If one becomes NULL, simply connect the other remaining

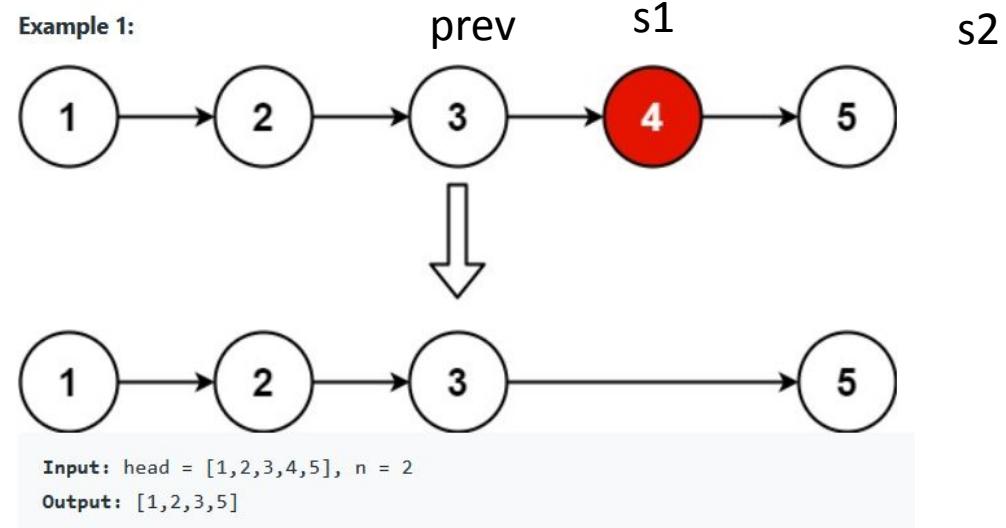
LC LinkedList#4: 23. Merge k Sorted Lists.(HARD)

```
11 class Cmp {
12     public:
13     bool operator()(ListNode* n1, ListNode* n2) {
14         return n1->val > n2->val; // ascending order
15     }
16 };
17 class Solution {
18 public:
19     static bool cmp(ListNode* n1, ListNode *n2) {
20         return n1->val < n2->val;
21     }
22
23     ListNode* mergeKLists(vector<ListNode*>& lists) {
24         priority_queue<ListNode *, vector<ListNode *>, Cmp> miniHeap;
25         for (int i = 0; i < lists.size(); i++) {
26             ListNode *head = lists[i];
27             while (head) {
28                 miniHeap.push(head);
29                 head = head->next;
30             }
31         }
32         ListNode *root = NULL, *cur = NULL, *tmp;
33         while (!miniHeap.empty()) {
34             ListNode *tmp = miniHeap.top();
35             tmp->next = NULL;
36             if (!root) {
37                 root = tmp;
38                 cur = root;
39             } else {
40                 cur->next = tmp;
41                 cur = tmp;
42             }
43             miniHeap.pop();
44         }
45     }
46     return root;
47 }
```

- Create a miniHeap using priority_queue (or simply use vector sorting)
- Traversal all K lists, add all nodes into minheap
- Pop minheap node one by one and form a new list to return

LC LinkedList#5: 19. Remove Nth Node From End of List.

```
-- 11  class Solution {
12  public:
13  ListNode* removeNthFromEnd(ListNode* head, int n) {
14      ListNode *s1 = head, *s2 = head, *prev = NULL;
15      for (int i = 0; i < n; i++) {
16          s2 = s2->next;
17      }
18      while (s1 && s2) {
19          prev = s1;
20          s1 = s1->next;
21          s2 = s2->next;
22      }
23      if (prev)
24          prev->next = s1->next;
25      else
26          return s1->next;
27
28      return head;
29  }
30 };
```



- Use two pointers: s2 is N step ahead of s1
- When s2 becomes NULL, S1 is the node to be removed
- Use “prev” hold node before s1

LC LinkedList#6: 143. Reorder List.

```
11 class Solution {
12 public:
13     void reorderList(ListNode* head) {
14         vector<ListNode*> nodes;
15         ListNode* cur = head;
16         while (cur) { // put all nodes into vector list with next set to NULL
17             ListNode *tmp = cur->next;
18             cur->next = NULL;
19             nodes.push_back(cur);
20             cur = tmp;
21         }
22         /////////////////////////////////
23         int N = nodes.size() -1;
24         cur = NULL;
25         for (int i = 0; i < N; i++) {
26             if (i < (N-i)) {
27                 nodes[i]->next = nodes[N-i];
28                 if (cur)
29                     cur->next = nodes[i];
30                 cur = nodes[N-i];
31             } else if (i == (N-i)) { // point to same node
32                 cur->next = nodes[i]; // last node
33             } else {
34                 break;
35             }
36         }
37     }
38 }
```

You are given the head of a singly linked-list. The list can be represented as:

$L_0 \rightarrow L_1 \rightarrow \dots \rightarrow L_{n-1} \rightarrow L_n$

Reorder the list to be on the following form:

$L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1} \rightarrow L_2 \rightarrow L_{n-2} \rightarrow \dots$

You may not modify the values in the list's nodes. Only nodes themselves may be changed.

- Push all ListNode into vector list
- Take out node from vector as required to form a new list.

m-by-n matrix

$$a_{i,j} \quad \text{n columns} \xrightarrow{\text{j changes}} \\ \begin{matrix} \text{m} \\ \text{rows} \end{matrix} \quad \begin{matrix} \downarrow \\ \text{i} \\ \text{c} \\ \text{h} \\ \text{a} \\ \text{g} \\ \text{e} \\ \text{s} \end{matrix} \quad \left[\begin{array}{cccc} a_{1,1} & a_{1,2} & a_{1,3} & \dots \\ a_{2,1} & a_{2,2} & a_{2,3} & \dots \\ a_{3,1} & a_{3,2} & a_{3,3} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{array} \right]$$

LC Matrix#1: 73. Set Matrix Zeroes.

```
1 class Solution {
2     public:
3         void setZeroes(vector<vector<int>>& matrix) {
4             int R = matrix.size();
5             int C = matrix[0].size();
6             unordered_set<int> r_zero;
7             unordered_set<int> c_zero;
8
9             for (int r = 0; r < R; r++) {
10                 vector<int> row = matrix[r];
11                 for (int c = 0; c < C; c++) {
12                     if (matrix[r][c] == 0) {
13                         if (r_zero.find(r) == r_zero.end()) {
14                             // set this row to zero
15                             r_zero.insert(r);
16                         }
17                         if (c_zero.find(c) == c_zero.end()) {
18                             // set this column to zero
19                             c_zero.insert(c);
20                         }
21                     }
22                 }
23             }
24             for (int r = 0; r < R; r++) {
25                 vector<int> row = matrix[r];
26                 for (int c = 0; c < C; c++) {
27                     if (r_zero.find(r) != r_zero.end()
28                         || c_zero.find(c) != c_zero.end()) {
29                         matrix[r][c] = 0;
30                     }
31                 }
32             }
33         }
34     };
```

Example 1:

1	1	1
1	0	1
1	1	1

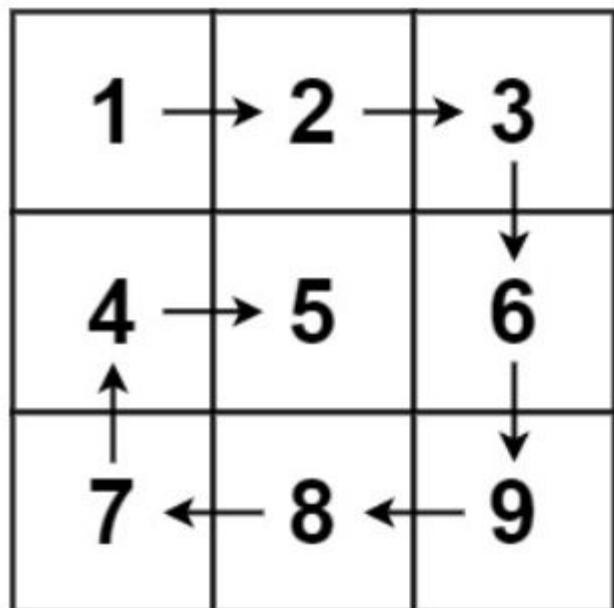
Input: matrix = [[1,1,1],[1,0,1],[1,1,1]]
Output: [[1,0,1],[0,0,0],[1,0,1]]

- Scan thru matrix to mark down rows and columns (into hashset) to be set to zero
- Second time go thru matrix again and check hashSets , either r or c in the its set, set matrix cell to be zero

LC Matrix#2: 54. Spiral Matrix.

```
1 class Solution {
2     public:
3         vector<int> spiralOrder(vector<vector<int>>& matrix) {
4             vector<int> res;
5             vector<vector<int>> DIRS {{0, 1}, {1, 0}, {0, -1}, {-1, 0}};
6             int R = matrix.size();
7             int C = matrix[0].size();
8             int r = 0, c = 0, mov_dir = 0;
9             int VISITED = INT_MAX;
10            while (1) {
11                int r_off = DIRS[mov_dir][0];
12                int c_off = DIRS[mov_dir][1];
13                res.push_back(matrix[r][c]);
14                matrix[r][c] = VISITED;
15                int nr = r + r_off;
16                int nc = c + c_off;
17                if (nr < 0 || nr >= R || nc < 0 || nc >= C || matrix[nr][nc] == INT_MAX) {
18                    // time to change direction
19                    mov_dir = (mov_dir + 1)%4;
20                    r = r+DIRS[mov_dir][0];
21                    c = c+DIRS[mov_dir][1];
22                    if (r < 0 || r >= R || c < 0 || c >= C) break;
23                    if (matrix[r][c] == INT_MAX) break;
24                } else {
25                    r = nr;
26                    c = nc;
27                }
28            }
29            return res;
30        }
31    };
```

Example 1:



Input: matrix = [[1,2,3],[4,5,6],[7,8,9]]

Output: [1,2,3,6,9,8,7,4,5]

- Set up four directions: L->R, U->D, R->L, D->U with offset
- Starts with [0][0], each visited node filled with INT_MAX
- When we move node, change direction when either reach out of boundary or VISTED node

LC Matrix#3: 48. Rotate Image.

You are given an $n \times n$ 2D matrix representing an image, rotate the image by 90 degrees (clockwise).

```
1 class Solution {
2     public:
3         void rotate(vector<vector<int>>& matrix) {
4             int N = matrix.size();
5             for (int r = 0; r < N; r++) {
6                 for (int c = r; c < N; c++) { // column starts with r !!!
7                     int t = matrix[r][c];
8                     matrix[r][c] = matrix[c][r];
9                     matrix[c][r] = t;
10                }
11            }
12
13            // roate clockwise , reverse each row
14            for (int r = 0; r < N; r++) {
15                int l_ = 0, r_ = N-1;
16                while (l_ < r_) {
17                    int t = matrix[r][l_];
18                    matrix[r][l_] = matrix[r][r_];
19                    matrix[r][r_] = t;
20
21                    l_++; r_--;
22                }
23            }
24 #ifdef ROTATE_COUNTERCLOCK
25             // if rotate counterlockwise, reverse each column
26             for (int c = 0; c < N; c++) {
27                 int t_ = 0, b_ = N-1;
28                 while (t_ < b_) {
29                     int t = matrix[t_][c];
30                     matrix[t_][c] = matrix[b_][c];
31                     matrix[b_][c] = t;
32
33                     t_++; b_--;
34                 }
35             }
36         #endif
37     }
38 }
```

Example 1:

1	2	3
4	5	6
7	8	9

Input: matrix = [[1,2,3],[4,5,6],[7,8,9]]
Output: [[7,4,1],[8,5,2],[9,6,3]]

- Flip value diagnose
- Reverse each row for clockwise
- (Reverse each column for counter-clockwise)

LC Matrix#4: 79. Word Search.

- Key points:

```
1 class Solution {
2     public:
3         int R = 0;
4         int C = 0;
5         //vector<vector<bool>> map;
6
7     bool backtrack(int r, int c, int pos, string word, vector<vector<char>>& board , vector<vector<bool>> map) {
8         if (pos == word.length()-1) return true;
9
10        if (map[r][c]) return false;
11
12        map[r][c] = true;
13        // check around neighbors (T, D, L, R) if it is matches next char, then call recursively
14        vector<vector<int>> DIRS { {-1, 0}, {1, 0}, {0,1}, {0,-1} };
15        for (int i = 0; i < DIRS.size(); i++) {
16            int nr = DIRS[i][0] + r;
17            int nc = DIRS[i][1] + c;
18            if (nr < 0 || nr >= R || nc < 0 || nc >= C || map[nr][nc]) continue;
19
20            if (board[nr][nc] == word.at(pos+1)) {
21                vector<vector<bool>> nmap = map; // pass new map down
22                if (backtrack(nr, nc, pos+1, word, board, nmap))
23                    return true;
24            }
25        }
26        return false;
27    }
28
29    bool exist(vector<vector<char>>& board, string word) {
30        int r = 0, c = 0, pos = 0;
31        R = board.size();
32        C = board[0].size();
33
34        vector<vector<bool>> map;
35        for (int i = 0; i < R; i++) {
36            vector<bool> t(C, false);
37            map.push_back(t);
38        }
39
40        // find all possible first character to get started
41        for (int r = 0; r < R; r++) {
42            for (int c = 0; c < C; c++) {
43                if (board[r][c] == word.at(0))
44                    if (backtrack(r, c, pos, word, board, map))
45                        return true;
46            }
47        }
48        return false;
49    }
};
```

Example 1:

A	B	C	E
S	F	C	S
A	D	E	E

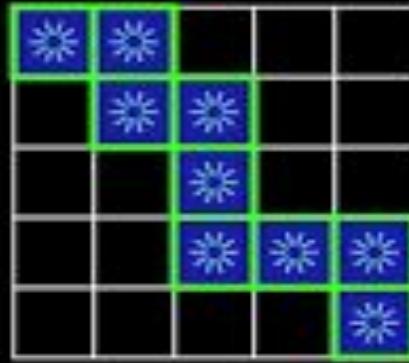
Input: board = [["A","B","C","E"],["S","F","C","S"],["A","D","E","E"]], word = "ABCED"

Output: true

LC Matrix#5: TBA.

- $(A \& B) << 1 \Rightarrow$ carry, $A^B \Rightarrow$ answer

Dynamic Programming



LC DP#1: 70. Climbing Stairs.

```
1  class Solution {
2  public:
3      vector<int> memo;
4      bool init = true;
5      int climbStairs(int n) {
6          if (n <=2) return n;
7          if (init) {
8              memo = vector<int>(n+1, 0);
9              init = false;
10         }
11         if (memo[n] > 0) return memo[n];
12         int res = climbStairs(n-1) + climbStairs(n-2);
13         memo[n] = res;
14         return res;
15     }
```

- $F(n) = F(n-1) + F(n-2)$ with $F(1) = 1$, $F(2) = 2$

LC DP#2: 322. Coin change.

```
3 class Solution {
4     public:
5         int coinChange(vector<int>& coins, int amount) {
6             vector<int> result(amount+1, INT_MAX);
7             result[0] = 0;
8             for (int i = 1; i <= amount; i++) {
9                 int ans = INT_MAX-1;
10                for (int j = 0; j < coins.size(); j++) {
11                    int index = i - coins[j];
12                    if (index >= 0) {
13                        ans = min(ans, 1 + result[index]);
14                    }
15                }
16                result[i] = ans;
17            }
18            return result[amount] >= INT_MAX-1? -1: result[amount];
19        }
20    }
```

Example 1:

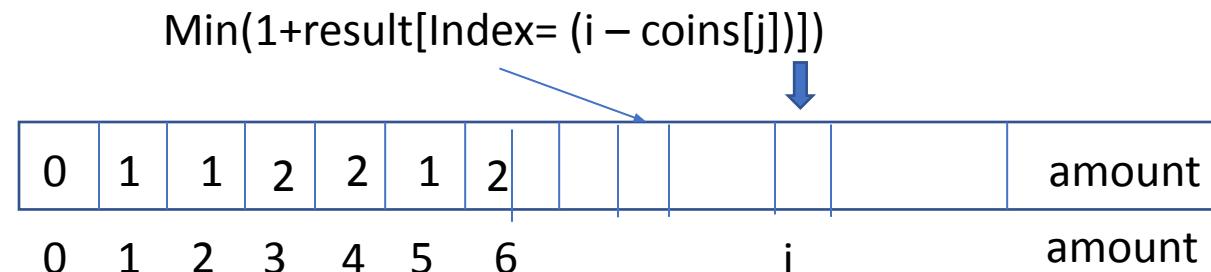
Input: coins = [1,2,5], amount = 11

Output: 3

Explanation: $11 = 5 + 5 + 1$

- for each amount, check for each coin
 - since coin is always positive, if subtract coin[i] is positive, we use it as index to get a result from that amount, plus 1 (current coin)
 - we do for all coins to find the mini number
 - Same nature of problem is perfect square LC#279, which is number of [1, 4, 9, 16...]

```
vector<int> result
```



LC DP#3: 300. Longest Increasing Subsequence(LIS) of given array.

Example 1:

```
1 class Solution {  
2     public:  
3         int lengthOfLIS(vector<int>& nums) {  
4             vector<int> res(nums.size(), 1);  
5  
6             for (int i = 1; i < nums.size(); i++) {  
7                 int ans = 1;  
8                 for (int j = 0; j <= i; j++) { // going thru all values before me,  
9                     if (nums[i] > nums[j]) { // for all values I'm greater  
10                         int t = res[j] + 1;  
11                         ans = max(ans, t); // keep max number  
12                     }  
13                 }  
14                 res[i] = ans;  
15             }  
16             int ret = 0;  
17             for (int i = 0; i < nums.size(); i++) ret = max(ret, res[i]);  
18             return ret;  
19         }  
20     };
```

Input: nums = [10,9,2,5,3,7,101,18]

Output: 4

Explanation: The longest increasing subsequence is [2,3,7,101], therefore the length is 4.

- Start from beginning, for each position, check all values before me
- If it is smaller than me, take its result +1
- Keep max value as result
- If it is Longest Decreasing Subsequence: starting from the end.

LC DP#4: 1143. Longest Common Sequence(LCS) of two strings.

```

1 class Solution {
2 public:
3     unordered_map<string, int> dp;
4
5 #ifdef RECURSION
6     int longestCommonSubsequence(string text1, string text2) {
7         if (text1.length() == 0 || text2.length() == 0) return 0;
8         if (dp.find(text1 + "-" + text2) != dp.end()) {
9             return dp[text1 + "-" + text2];
10        }
11        int res = 0;
12
13        string t1 = text1.substr(1);
14        string t2 = text2.substr(1);
15        if (text1.at(0) == text2.at(0)) {
16            res = 1 + longestCommonSubsequence(t1, t2);
17        } else {
18            int lcs1 = longestCommonSubsequence(text1, t2);
19            int lcs2 = longestCommonSubsequence(t1, text2);
20            res = max(lcs1, lcs2);
21        }
22        dp[text1 + "-" + text2] = res;
23        return res;
24    }
25 #endif
26     int longestCommonSubsequence(string text1, string text2) {
27         vector<vector<int>> dp(text1.length() + 1, vector<int>(text2.length() + 1, 0));
28         int r = text1.length() - 1;
29         int c = text2.length() - 1;
30
31         while (r >= 0 && c >= 0) {
32             for (int c_ = c; c_ >= 0; c--) {
33                 if (text2.at(c_) == text1.at(r))
34                     dp[r][c_] = 1 + dp[r+1][c_+1];
35                 else
36                     dp[r][c_] = max(dp[r+1][c_], dp[r][c_+1]);
37             }
38             for (int r_ = r; r_ >= 0; r--) {
39                 if (text2.at(c) == text1.at(r_))
40                     dp[r_][c] = 1 + dp[r_+1][c];
41                 else
42                     dp[r_][c] = max(dp[r_+1][c], dp[r_][c+1]);
43             }
44             r--;
45             c--;
46         }
47         return dp[0][0];
48     }
};

```

- Use recursion: straightforward, but it takes much more time $O(2^{\text{power min}(N, M)})$
- Basically below plus memo:
 $\text{if } (\text{text1.at}(0) == \text{text2.at}(0))$
 $\quad \text{return } 1 + \text{LCS}(\text{text1.substr}(1), \text{text2.substr}(1));$
 else
 $\quad \text{return } (\max(\text{LCS}(\text{text1}, \text{text2.substr}(1)),$
 $\quad \quad \text{LCS}(\text{text1.substr}(1), \text{text2})));$
- Use dp approach: it is $O(m * n)$

A	B	C	K	O	K	G	E
1	0						
0	1	0					
	1	0					
		1	0				
		0	1	0			
			1	0			
			0	1	0		
				0	1	0	
					0	1	0
						0	1
							0

LC DP#5: 139. word break: given sentence and word dictionary, tell if it can be break using word from dictionary.

```
1 class Solution {
2 public:
3     unordered_map<string, bool> dp;
4
5     bool wordBreak(string s, vector<string>& wordDict) {
6         if (s.length() == 0) return true;
7         if (dp.find(s) != dp.end()) return dp[s];
8
9         // check if any matched word in dict,
10        for (int i = 0; i < wordDict.size(); i++) {
11            string word=wordDict[i];
12            if (s.find(word, 0) == 0) {
13                if(wordBreak(s.substr(word.length()), wordDict)) {
14                    dp[s] = true;
15                    return true;
16                }
17            }
18        }
19        dp[s] = false;
20        return false;
21    }
22};
```

- std::string::find(substr, pos): find first occurrence substr starting index of pos
- std::string::rfind(substr, pos): find last occurrence substr starting from pos
- std::string::find_first_of(char c, int pos)
- Go thru given dictionary, for each word, use s.find(word, 0) to see if given string starts with word, if yes, recursive call of substring (after remove word)
- Use memo to improve performance

Backtracking recipe

```
void Backtrack(res, args)
    if ( GOAL REACHED )
        add solution to res
        return
    for ( int i = 0; i < NB_CHOICES; i++ )
        if ( CHOICES[i] is valid )
            make choices[i]
            Backtrack(res, args)
            undo choices[i]
```

LC DP#6.0 permutation

```
3     // https://www.youtube.com/watch?v=Nabbpl7y4Lo
4     // three steps of backtrack:
5     // #1 check if reaches goal, if yes, save result and return
6     // #2 among all possible choices
7     //   - take valid choice, update in input list
8     //   -- perform bacKtrack
9     //   - undo the choice, go for next choice
10    //
11    void backtrack(vector<int>& nums, vector<int>& perm, vector<vector<int>>& res, vector<bool>& used) {
12        if (perm.size() == nums.size()) { // check if meets goal
13            res.push_back(perm);           // yes, save result
14            return;
15        }
16        for (int i = 0; i < nums.size(); i++) { // all possible choices
17            if (!used[i]) {                // only take valid one
18                used[i] = true;           // make the choice
19                perm.push_back(nums[i]); // update parameters
20                backtrack(nums, perm, res, used);
21                perm.pop_back();         // undo the choice
22                used[i] = false;          // undo the |choice
23            }
24        }
25    }
26 }
27 vector<vector<int>> permute(vector<int>& nums) {
28     vector<vector<int>> res;
29     vector<int> perm;
30     vector<bool> used(nums.size(), false);
31     backtrack(nums, perm, res, used);
32     return res;
33 }
34 }
```

LC DP#6: 39. combination sum I

Given an array of **distinct** integers `candidates` and a target integer `target`, return a *list of all unique combinations* of `candidates` where the chosen numbers sum to `target`. You may return the combinations in **any order**.

The **same** number may be chosen from `candidates` an **unlimited number of times**. Two combinations are unique if the frequency of at least one of the chosen numbers is different.

It is **guaranteed** that the number of unique combinations that sum up to `target` is less than `150` combinations for the given input.

```
1  class Solution {
2  public:
3      // reference: https://www.youtube.com/watch?v=yFfv03AE_vA
4      void backTrack(vector<int>& candidates, int start, vector<int> list, vector<vector<int>>& result, int target) {
5          if (target < 0) return;
6          if (target == 0) {
7              if (list.size() > 0) {
8                  vector<int> res1(list);
9                  result.push_back(res1);
10             }
11             return;
12         }
13         for (int i = start; i < candidates.size(); i++) { // start is used to prevent duplicate
14             list.push_back(candidates[i]);
15             backTrack(candidates, i, list, result, target-candidates[i]);
16             list.pop_back();
17         }
18     }
19
20     vector<vector<int>> combinationSum(vector<int>& candidates, int target) {
21         vector<vector<int>> result;
22         vector<int> list;
23
24         backTrack(candidates, 0, list, result, target);
25
26         return result;
27     }
28 }
```

Allow same number to be re-used

- Numbers are distinct
- backtrack allow same number to be re-used

LC DP#6: 40. combination sum II

Given a collection of candidate numbers (`candidates`) and a target number (`target`), find all unique combinations in `candidates` where the candidate numbers sum to `target`.

Each number in `candidates` may only be used **once** in the combination.

```
1 class Solution {
2     public:
3         void backTrack(vector<int>& candidates, int start, int target, vector<int> list, vector<vector<int>>& result) {
4             if (target < 0) return;
5             if (target == 0) {
6                 vector<int> ans(list);
7                 result.push_back(ans);
8                 return;
9             }
10            for (int i = start; i < candidates.size(); i++) {
11                // this makes sure number is used only once
12                if (i > start && candidates[i] == candidates[i-1]) continue; ←
13                if (target - candidates[i] < 0) break; // since it is sorted
14                list.push_back(candidates[i]);
15                backTrack(candidates, i+1, target - candidates[i], list, result); // i+1 to avoid re-use the same number
16                list.pop_back();
17            }
18        }
19
20        vector<vector<int>> combinationSum2(vector<int>& candidates, int target) {
21            vector<vector<int>> result;
22            vector<int> list;
23            sort(candidates.begin(), candidates.end()); // sort in ascending order
24            backTrack(candidates, 0, target, list, result);
25            return result;
26        }
27    };
```

- Values are random, could duplicate
- Sort them in ascending order
- each occurrence of number is used only once
- **Skip the following numbers with same value**

LC DP#6: 216. combination sum III

Find all valid combinations of `k` numbers that sum up to `n` such that the following conditions are true:

- Only numbers `1` through `9` are used.
- Each number is used **at most once**.

Return *a list of all possible valid combinations*. The list must not contain the same combination twice, and the combinations may be returned in any order.

```
1 class Solution {
2 public:
3     void backtrack(int k, int start, int target, vector<int> list, vector<vector<int>> &result) {
4         if (target < 0) return;
5         if (target == 0) {
6             if (list.size() == k) { // result only allow K numbers
7                 vector<int> ans(list);
8                 result.push_back(ans);
9             }
10            return;
11        }
12        if (list.size() >= k) return; // too many stop here
13        for (int i = start; i < 10; i++) { // only use the digits 1-9
14            list.push_back(i);
15            backtrack(k, i+1, target-i, list, result);
16            list.pop_back();
17        }
18    }
19
20    vector<vector<int>> combinationSum3(int k, int n) {
21        vector<vector<int>> result;
22        vector<int> list;
23        backtrack(k, 1, n, list, result);
24        return result;
25    }
26};
```

- simplified version of II

LC DP#6: 377. combination sum IV

Given an array of **distinct** integers `nums` and a target integer `target`, return *the number of possible combinations that add up to target*.

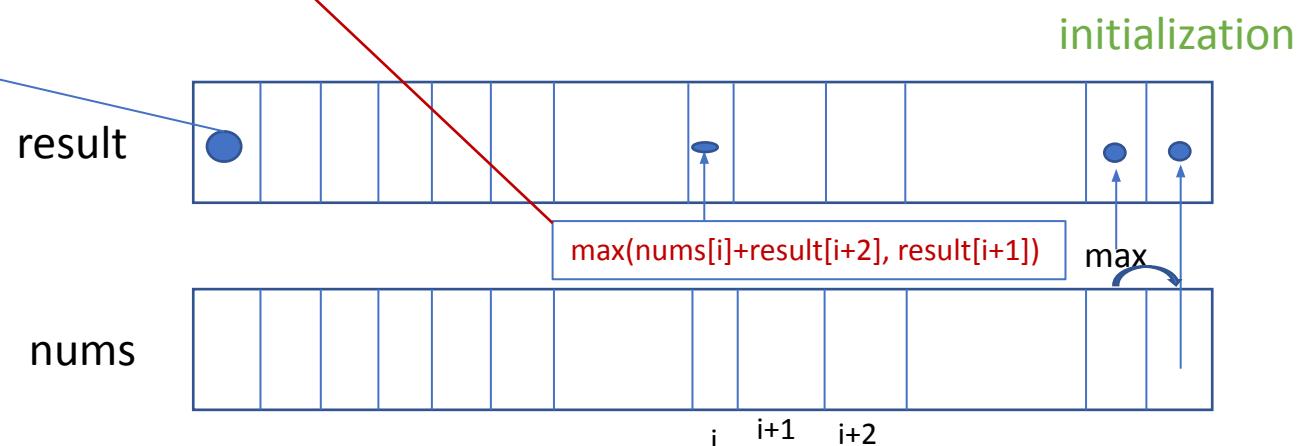
```
1 class Solution {
2     public:
3         unordered_map<string, int> memo;
4         int backtrack(vector<int>& nums, int pos, int target) {
5             int res = 0;
6             if (target < 0) return 0;
7             if (target == 0) {
8                 return 1;
9             }
10            string key = to_string(pos) + "----" + to_string(target);
11            if (memo.find(key) != memo.end()) return memo[key];
12
13            for (int i = pos; i < nums.size(); i++) {
14                if (target - nums[i] < 0) break;
15                res += backtrack(nums, pos, target - nums[i]);
16            }
17            memo[key] = res;
18            return res;
19        }
20
21        int combinationSum4(vector<int>& nums, int target) {
22            int result = 0;
23            sort(nums.begin(), nums.end());
24            result = backtrack(nums, 0, target);
25            return result;
26        }
27    };
```

- backtrack with memo of pos+target

LC DP#7: 198. Rob house: Given an integer array nums representing the amount of money of each house, return the maximum amount of money you can rob tonight without alerting the police.

```
1 class Solution {
2     public:
3         int rob(vector<int>& nums) {
4             if (nums.size() == 1) return nums[0];
5
6             vector<int> result(nums.size(), 0);
7             result[nums.size()-1] = nums[nums.size()-1];
8             result[nums.size()-2] = max(nums[nums.size()-1], nums[nums.size()-2]) ;
9             for (int i = nums.size()-3; i >= 0; i--) {
10                 result[i] = max(result[i+2]+nums[i], result[i+1]);
11             }
12             return result[0];
13         }
14     };
15 }
```

- Since we don't know what is ahead to make it maximum, we start from the end: last one is it self, last second one is if it is greater last, then it is self, otherwise would be last one
- For any given house position "index", compare "its value + next_next" with "next", take the bigger value



LC DP#8: 91. decode way. 226 => BBF, BZ, VF

A message containing letters from `A-Z` can be **encoded** into numbers using the following mapping:

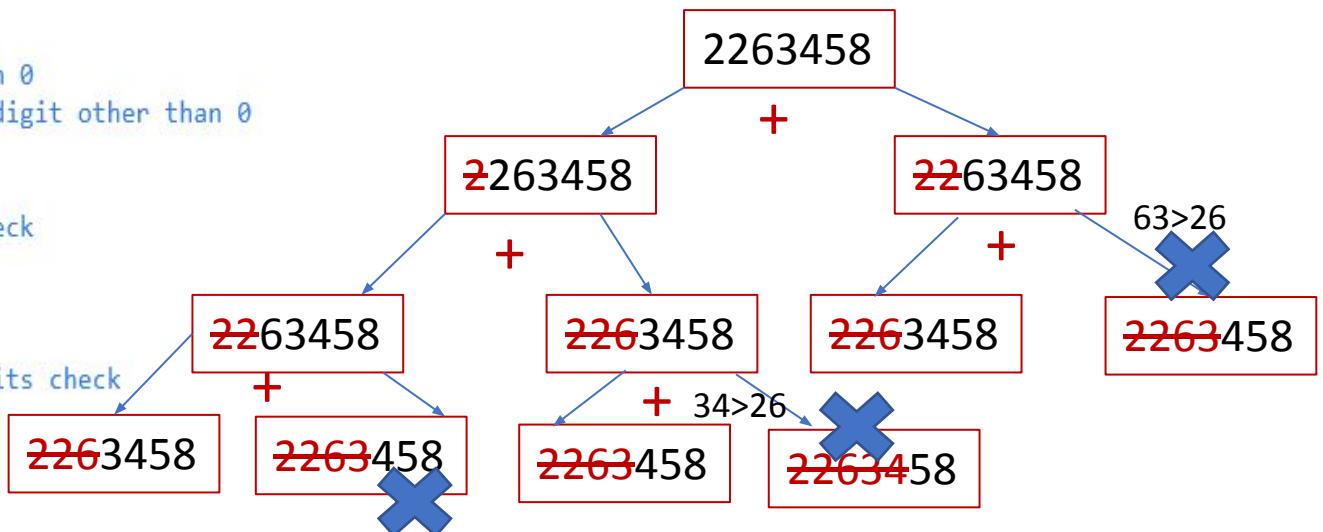
```
'A' -> "1"  
'B' -> "2"  
...  
'Z' -> "26"
```

To **decode** an encoded message, all the digits must be grouped then mapped back into letters using the reverse of the mapping above (there may be multiple ways). For example, `"11106"` can be mapped into:

- "AAJF" with the grouping `(1 1 10 6)`
- "KJF" with the grouping `(11 10 6)`

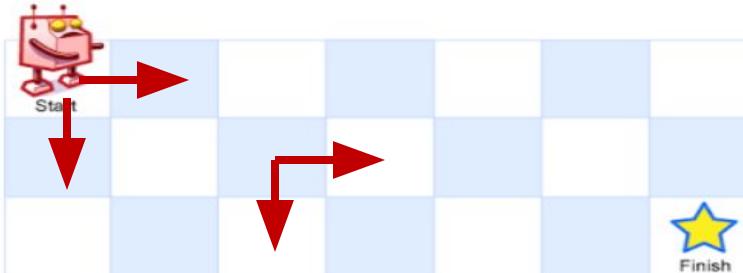
```
3 class Solution {  
4     private :  
5     public:  
6         unordered_map<string, int> memo;  
7  
8     int numDecodings(string s) {  
9         if (s.length() == 0) return 1;  
10        if (s.at(0) == '0') return 0; // can't decode starting with 0  
11        if (s.length() == 1) return 1; // good if it is only one digit other than 0  
12        if (memo.find(s) != memo.end()) return memo[s];  
13  
14        int res = numDecodings(s.substr(1)); // spin one digit check  
15        if (s.length() >= 2) {  
16            int val = stoi(s.substr(0,2));  
17            if (val <= 26) {  
18                res += numDecodings(s.substr(2)); // spin two digits check  
19            }  
20        }  
21        memo[s] = res;  
22        return res;  
23    }  
24};
```

- Check special cases:
 - end of string: good (return 1)
 - Starts with 0, can't decode
 - One digit other than 0, good (return 1)
- Check memo cache, if it exists, return result
- Otherwise try to explore with one digit and two digits if possible



LC DP#9: 62. unique paths/grid walk

Example 1:



Input: m = 3, n = 7

Output: 28

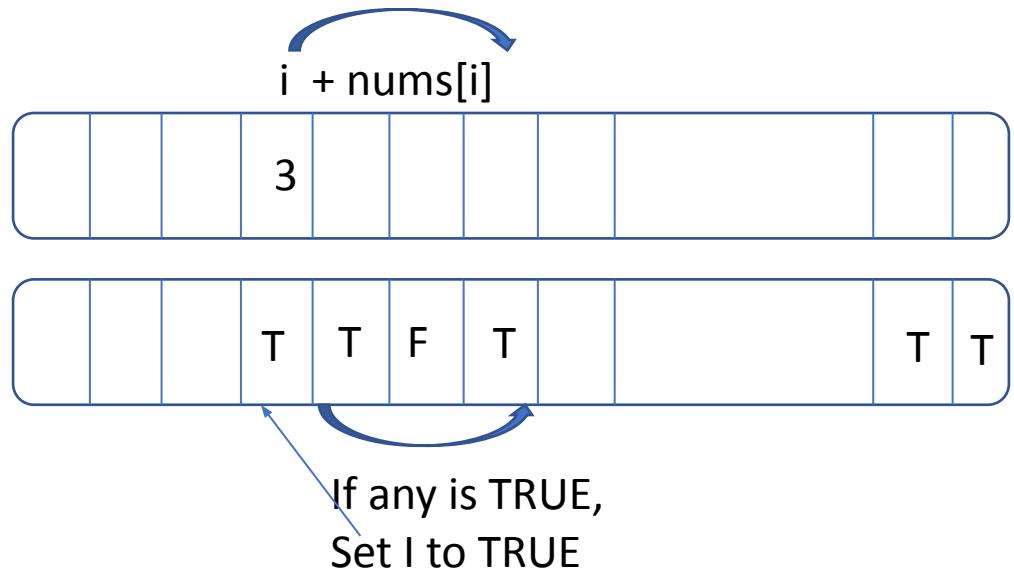
- $\text{Res}[m,n] = \text{res}[m-1, n] + \text{res}[m, n-1]$
- Use a `unordered_map<string, int>` as memo
- Key is “ “r”==“c”

```
1 class Solution {
2     private:
3         unordered_map<string, int> memo;
4     public:
5
6     int uniquePaths(int m, int n) {
7
8         if (m < 1 || n < 1) return 0;
9         if (m == 1 || n == 1) return 1; // walk bottom or right
10        string key = to_string(m) + " --- " + to_string(n);
11        if (memo.find(key) != memo.end()) {
12            return memo[key];
13        }
14        int res = uniquePaths(m-1, n) + uniquePaths(m, n-1);
15        memo[key] = res;
16
17        return res;
18    }
19};
```

LC DP#10: 55. Jump Game.

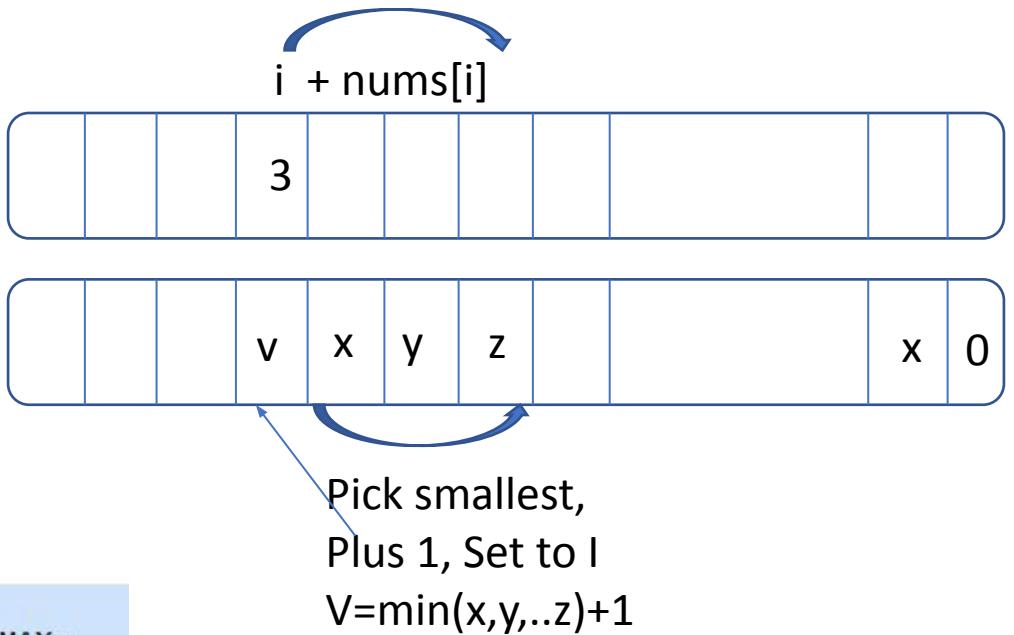
```
class Solution {
public:
    bool canJump(vector<int>& nums) {
        int N = nums.size();
        vector<bool> status(N, false);

        status[N-1] = true; // last stop is true
        for (int i = N - 2; i >= 0; i--) {
            int farestJump = min(i + nums[i], N-1);
            for (int j = i+1; j <= farestJump; j++) {
                if (status[j] == true) { // within reach, if any is good
                    status[i] = true; // we are good.
                    break;
                }
            }
        }
        return status[0];
    }
};
```



LC DP#11: 45. Jump Game II.

```
1 class Solution {
2 public:
3     int jump(vector<int>& nums) {
4         int N = nums.size();
5         vector<int> jumps(N);
6         jumps[N-1] = 0;
7
8         for (int i = N-2; i >= 0; i--) {
9             int maxPos = min(N-1, i+nums[i]);
10            int ans = INT_MAX;
11            for (int j=i+1; j <= maxPos; j++) {
12                int tmp = jumps[j] != INT_MAX ? (jumps[j] + 1):INT_MAX;
13                ans = min(ans, tmp);
14            }
15            jumps[i] = ans;
16        }
17        return jumps[0];
18    }
19};
```



LC DP#12: 1306. Jump Game III

Given an array of non-negative integers `arr`, you are initially positioned at `start` index of the array. When you are at index `i`, you can jump to `i + arr[i]` or `i - arr[i]`, check if you can reach to **any** index with value 0.

```
1 class Solution {
2 public:
3     bool helper(vector<int>& arr, int start, vector<bool>& v) {
4         if (arr[start] == 0) return true; // reach desired place
5         if (v[start] == true) return false; // found loop
6
7         v[start] = true; // mark this is visited
8         vector<bool> nv = v;
9         if (start + arr[start] < arr.size()) { // try next round
10             if (helper(arr, start + arr[start], nv)) return true;
11         }
12         if (start - arr[start] >= 0) {
13             if (helper(arr, start - arr[start], nv)) return true;
14         }
15         return false;
16     }
17
18     bool canReach(vector<int>& arr, int start) {
19         vector<bool> v(arr.size(), false);
20
21         return helper(arr, start, v);
22     }
23 };
```

- Recursive call forward and backward if it inbound
- Use a visited map, if reach visited pos, loop detected, NOT possible

Example 1:

Input: arr = [4,2,3,0,3,1,2], start = 5

Output: true

Explanation:

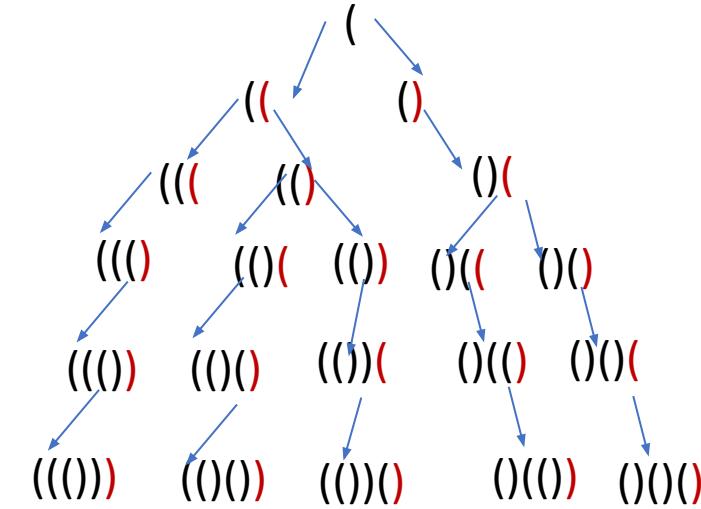
All possible ways to reach at index 3 with value 0 are:

index 5 → index 4 → index 1 → index 3

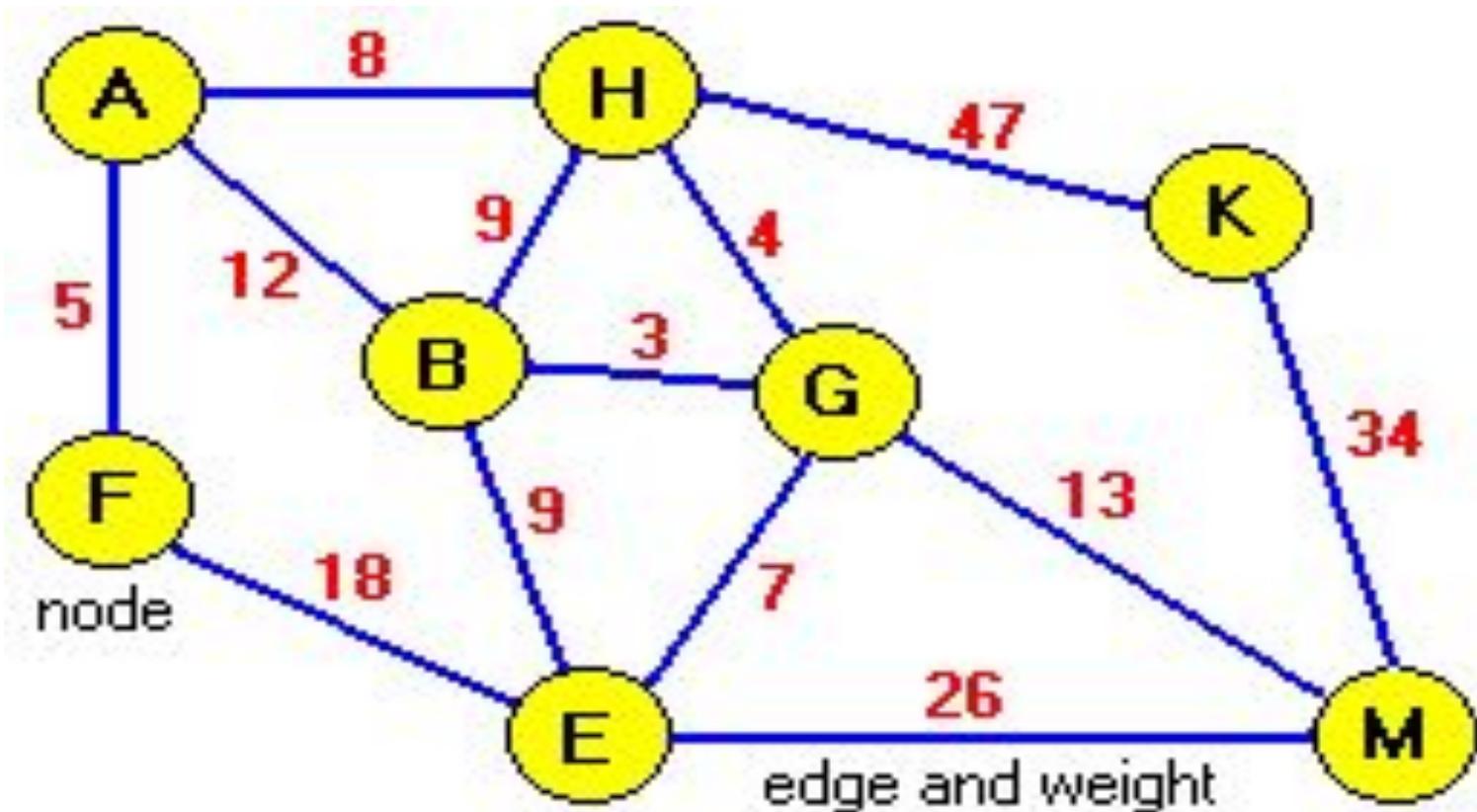
index 5 → index 6 → index 4 → index 1 → index 3

LC DP#13: 22: Generate parentheses

```
1 class Solution {
2 public:
3     int N;
4     void backTrack(vector<string>& res, vector<char> cur, int open, int close) {
5         if (cur.size() == 2*N) {
6             string ans = "";
7             for (int i = 0; i < 2*N; i++) ans.append(1,cur[i]);
8             res.push_back(ans);
9             return;
10        }
11        if (open < N) {
12            cur.push_back('(');
13            backTrack(res, cur, open+1, close);
14            cur.pop_back();
15        }
16        if (open > close) {
17            cur.push_back(')');
18            backTrack(res, cur, open, close+1);
19            cur.pop_back();
20        }
21    }
22    vector<string> generateParenthesis(int n) {
23        vector<string> res;
24        vector<char> cur;
25        N = n;
26        backTrack(res, cur, 0, 0);
27        return res;
28    }
29};
```



Graph



LC Graph#1: 133. clone graph.

```
22 class Solution {
23     public:
24         Node* cloneGraph(Node* node) {
25             Node *root = NULL;
26             queue<Node *> q;
27             unordered_map<int, Node *> newNodes;
28             unordered_set<int> seen;
29             if (node == NULL) return NULL;
30
31             q.push(node);
32
33             while (!q.empty()) {
34                 Node *t = q.front();
35                 q.pop();
36                 Node *n = NULL;
37                 if (newNodes.find(t->val) == newNodes.end()) {
38                     n = new Node(t->val);
39                     if (root == NULL) root = n;
40                     newNodes[t->val] = n;
41                 } else {
42                     n = newNodes[t->val];
43                 }
44
45                 for (int i = 0; i < t->neighbors.size(); i++) {
46                     Node *nei = t->neighbors[i];
47                     Node *nnei = NULL;
48                     if (newNodes.find(nei->val) == newNodes.end()) {
49                         nnei = new Node(nei->val);
50                         newNodes[nnei->val] = nnei;
51                     } else {
52                         nnei = newNodes[nei->val];
53                     }
54                     if (find(n->neighbors.begin(), n->neighbors.end(), nnei) == n->neighbors.end())
55                         n->neighbors.push_back(nnei);
56                     if (seen.find(nei->val) == seen.end())
57                         q.push(nei);
58                 }
59                 seen.insert(n->val);
60             }
61         }
62     };
63 }
```

```
3 class Node {
4     public:
5         int val;
6         vector<Node*> neighbors;
7         Node() {
8             val = 0;
9             neighbors = vector<Node*>();
10    }
11    Node(int _val) {
12        val = _val;
13        neighbors = vector<Node*>();
14    }
15 }
```

- Use BFS (i.e. use Q) to traversal Graph with help of `unordered_set<int>`
- Create a `unordered_map<int, Node*>` to track all created cloned Node, key is node#
- When traverse a node,
 - check cloned node in map, if it doesn't exist, create it.
 - find its all neighbor nodes, check if its cloned node is in map, if not created,
 - Add all cloned neighbors under cloned node

LC Graph#2: 207. Course Schedule.

There are a total of `numCourses` courses you have to take, labeled from `0` to `numCourses - 1`. You are given an array `prerequisites` where `prerequisites[i] = [ai, bi]` indicates that you **must** take course `bi` first if you want to take course `ai`. Return `true` if you can finish all courses. Otherwise, return `false`.

```
1  class Solution {
2      const int INIT = 0;
3      const int DONE = 1;
4      const int PROCESSING = 2;
5      unordered_map<int, vector<int>> preq_list; // essentially adjacency list of graph
6      vector<int> state;
7
8  public:
9      bool hasCycle(int course) {
10         if (state[course] == PROCESSING) return true; // cycle detected!
11         state[course] = PROCESSING;
12         if (preq_list.find(course) != preq_list.end()) {
13             vector<int> preqs = preq_list[course];
14             for (int j = 0; j < preqs.size(); j++) {
15                 if (state[preqs[j]] != DONE) {
16                     if (hasCycle(preqs[j]))
17                         return true;
18                 }
19             }
20             state[course] = DONE;
21             return false;
22         }
23     bool canFinish(int numCourses, vector<vector<int>>& prerequisites) {
24         // go thru to establish preq for each course (node)
25         for (int i = 0; i < prerequisites.size(); i++) {
26             vector<int> elem = prerequisites[i];
27             int course = elem[0];
28             int pre = elem[1];
29             if (preq_list.find(course) != preq_list.end()) {
30                 vector<int> courses = preq_list[course];
31                 courses.push_back(pre);
32                 preq_list[course] = courses; // update
33             } else {
34                 vector<int> preq { pre }; //preq.push_back(pre);
35                 preq_list[course] = preq;
36             }
37         }
38         // now go thru each course, to make sure there is no loop in the preq_list
39         state = vector<int>(numCourses, INIT); // essentially it creates an array, init to zero
40         for (int i = 0; i < numCourses; i++) {
41             if (state[i] == INIT) { // if this course was never initiated, start detect
42                 if (hasCycle(i))
43                     return false;
44             }
45         }
46     }
47     return true;
48 }
```

- Go thru given pre-requisites to establish `preq_list` `unordered_map<int, vector<int>>`: key is course#, value is its pre_requisite `vector<int>`
- Init state for each course to INIT
- Go thru each course, using GRAPH COLOR to detect if there is any cycle.

LC Graph#3: 210. Course Schedule II (topological sort).

```
1 class Solution {
2 public:
3     vector<int> findOrder(int numCourses, vector<vector<int>>& prerequisites) {
4         vector<int> res;
5         vector<vector<int>> preq_list(numCourses);
6         // establish indegree for each course: number of pre-requisites
7         vector<int> indegrees(numCourses, 0);
8         for (int i = 0; i < prerequisites.size(); i++) {
9             int target = prerequisites[i][0];
10            int condition = prerequisites[i][1];
11            preq_list[condition].push_back(target); // collect who requires/needs me
12            indegrees[target] += 1; // indegree table for each course
13        }
14
15        // go thru indegrees table to find initial courses with zero indegree
16        queue<int> q;
17        unordered_set<int> visited;
18        for (int i = 0; i < numCourses; i++) {
19            if (indegrees[i] == 0) {
20                q.push(i);
21                visited.insert(i);
22            }
23        }
24        while (!q.empty()) {
25            int c = q.front(); q.pop();
26            res.push_back(c);
27
28            // find who needs me, decrease indegrees by 1 ,
29            vector<int> targets = preq_list[c];
30            for (int j = 0; j < targets.size(); j++) {
31                indegrees[targets[j]] -= 1;
32            }
33            // and add the member with zero into q if it is not seen before
34            for (int i = 0; i < numCourses; i++) {
35                if (indegrees[i] == 0 && visited.find(i) == visited.end()) {
36                    q.push(i);
37                    visited.insert(i);
38                }
39            }
40        }
41        if (res.size() != numCourses) res.clear();
42        return res;
43    }
44};
```

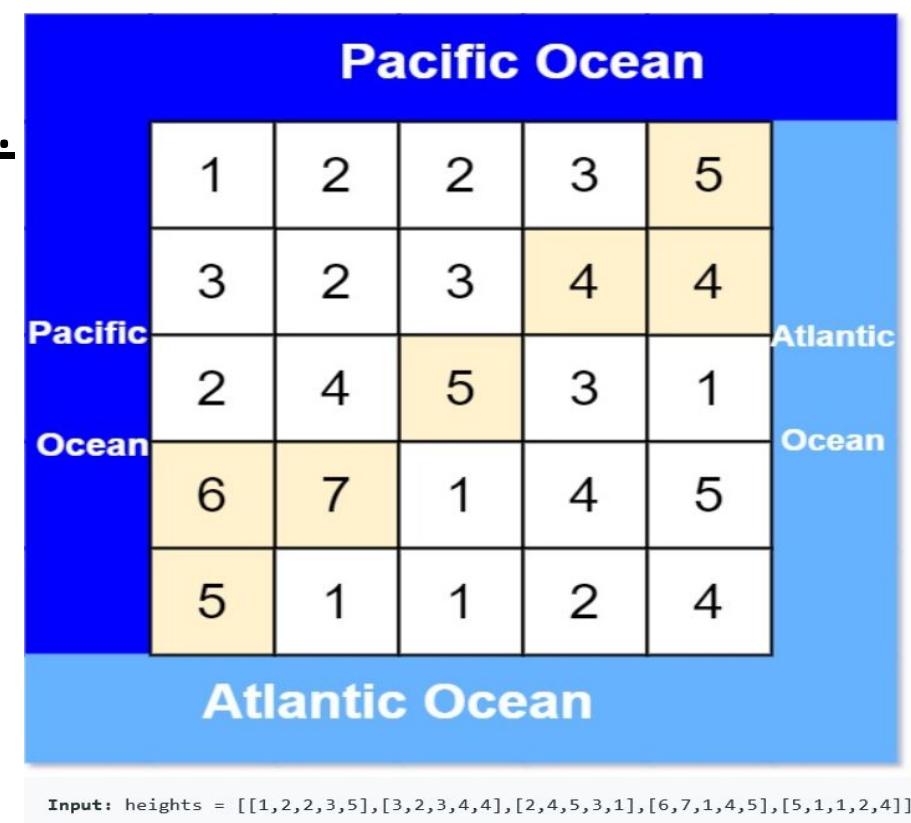
- Go thru given pre-requisites list to establish `indegree[target]` table and `preq_list<cond, targets>` (for a given course, all other courses require/need it)
- Add all course with 0 in indegree table into Q
- For each course in Q,
 - dequeue it,
 - Add it into result list
 - decrease 1 in indegree table for all other courses which require it.
 - Add any new courses with 0 indegree.
- Repeat until Q is empty

LC Graph#4: 417. Pacific Atlantic Water Flow.

```

1 class Solution {
2     const int PAC = 0x1;
3     const int ATL = 0x2;
4
5     public:
6         void dfs(int r, int c, vector<vector<int>>& h, int flag, vector<vector<int>>& map) {
7             vector<vector<int>> DIRS;
8             vector<int> up {-1, 0}, down {1,0}, left{0,-1}, right{0,1};
9             DIRS.push_back(up); DIRS.push_back(down); DIRS.push_back(left); DIRS.push_back(right);
10            int R = h.size();
11            int C = h[0].size();
12            map[r][c] |= flag;
13            // DFS explore thru 4 directions
14            for (auto d: DIRS) {
15                int nr = r + d[0];
16                int nc = c + d[1];
17                if (nr < 0 || nr >= R || nc < 0 || nc >= C ) continue; // out of space
18                if ((map[nr][nc] & flag) == flag) continue; // reached already
19                if (h[nr][nc] >= h[r][c]) dfs(nr, nc, h, flag, map); // if neighbor is not lower, explore
20            }
21        }
22        vector<vector<int>> pacificAtlantic(vector<vector<int>>& heights) {
23            int R = heights.size();
24            int C = heights[0].size();
25            vector<vector<int>> map(R, vector<int>(C,0));
26
27            // do DFS traversal starting from known cells:
28            // pacific: top and left
29            // atlantic: bottow and right
30            for (int r = 0; r < R; r++) {
31                dfs(r, 0, heights, PAC, map); // left side
32                dfs(r, C-1, heights, ATL, map); // right side
33            }
34            for (int c = 0; c < C; c++) {
35                dfs(0, c, heights, PAC, map); // top side
36                dfs(R-1, c, heights, ATL, map); // bottom side
37            }
38            vector<vector<int>> res;
39            for (int r = 0; r < R; r++) {
40                for (int c = 0; c < C; c++) {
41                    if (map[r][c] == (PAC|ATL)) {
42                        vector<int> ans {r,c};
43                        res.push_back(ans);
44                    }
45                }
46            }
47        }
48    }
49}

```



Input: heights = [[1,2,2,3,5],[3,2,3,4,4],[2,4,5,3,1],[6,7,1,4,5],[5,1,1,2,4]]
Output: [[0,4],[1,3],[1,4],[2,2],[3,0],[3,1],[4,0]]

- Use DFS find all cells (4 directions) which have equal or higher HEIGHT
- Pacific starts top & left cells; Atlantic starts bottom and right cells
- Use a vector<vector<int>> map to record DFS result
- Check each cell result in map, if both PAC and ATL set, add into result.

LC Graph#5: 200. Number of Islands.

```
1 class Solution {
2 public:
3     int R, C;
4     void dfs(int r, int c, vector<vector<char>>& grid) {
5         vector<vector<int>> DIRS { {-1,0}, {1,0}, {0,-1}, {0,1} };
6         grid[r][c] = 0; //visit and zero out
7         for (auto d: DIRS) {
8             int nr = r + d[0];
9             int nc = c + d[1];
10            if (nr < 0 || nc < 0 || nr >= R || nc >= C) continue;
11            if (grid[nr][nc] == '1') dfs(nr, nc, grid);
12        }
13    }
14
15    int numIslands(vector<vector<char>>& grid) {
16        R = grid.size();
17        C = grid[0].size();
18        int res = 0;
19
20        for (int r = 0; r < R; r++) {
21            for (int c = 0; c < C; c++) {
22                if (grid[r][c]=='1') {
23                    res++;
24                    dfs(r, c, grid);
25                }
26            }
27        }
28        return res;
29    }
30};
```

- Search thru grid, find 1,
- Increment result
- Perform DFS to zero out all neighbors(UP, BOTTOM, LEFT, RIGHT)
- Repeat above

LC Graph#6: 128. Longest Consecutive Sequence. Must be O(n)

```
1 class Solution {
2 public:
3     int longestConsecutive(vector<int>& nums) {
4         if (nums.size() == 0) return 0;
5         unordered_set<int> allNums(nums.begin(), nums.end());
6         unordered_map<int, bool> checked; // to ensure O(n) check
7         int res = 1;
8         for (int i = 1; i < nums.size(); i++) {
9             int n = nums[i];
10            if (checked.find(n) == checked.end()) {
11                int counter = 1;
12                int target = n - 1; // search downwards
13                while (allNums.find(target) != allNums.end()) {
14                    counter++;
15                    checked[target] = true;
16                    target--;
17                }
18                target = n + 1; // search upwards
19                while (allNums.find(target) != allNums.end()) {
20                    counter++;
21                    checked[target] = true;
22                    target++;
23                }
24                res = max(counter, res);
25                checked[n] = true;
26            }
27        }
28    }
29 }
30
```

- Add all numbers into a set
- Go thru each number, search downward and upwards consecutively
- Counter all of them
- Mark them checked in HashMap
- Return max counter

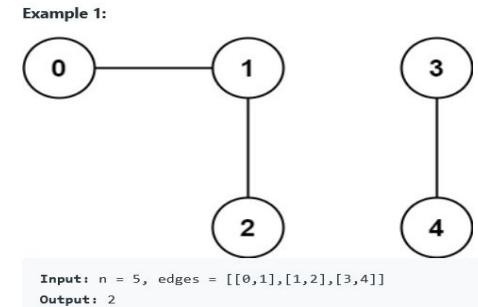
LC Graph#7: 261. Graph Valid Tree: i.e. all connected without loop.

```
1 class Solution {
2 public:
3     bool validTree(int n, vector<vector<int>>& edges) {
4         // establish adjList
5         vector<vector<int>> adj(n);
6         for (int i = 0; i < edges.size(); i++) {
7             int a = edges[i][0]; // use a, b make more sense since this is undirected graph
8             int b = edges[i][1];
9             adj[a].push_back(b); // every edge, we have two edges save, need remove
10            adj[b].push_back(a); // one edge during traverse the first one
11        }
12        unordered_set<int> seen; // save the traversed nodes as result
13        stack<int> st;
14        st.push(0); // start with first node
15        seen.insert(0);
16        while (!st.empty()) {
17            int n = st.top(); st.pop();
18            for (auto nei: adj[n]) {
19                if (seen.find(nei) != seen.end()) {
20                    return false; // there is cycle
21                }
22                seen.insert(nei);
23                st.push(nei);
24
25                // remove edge from nei->n since this is undirected graph
26                vector<int> nodes_2_nei = adj[nei];
27                nodes_2_nei.erase(find(nodes_2_nei.begin(), nodes_2_nei.end(), n));
28                adj[nei] = nodes_2_nei; // need update back!!!
29            }
30        }
31
32        return seen.size() == n;
33    }
34};
```

- A graph is a valid tree: all nodes are connected without loop
- Establish adjacency list for each node by going thru edge list: by directions
- DFS using STACK: Starts with first node, use a “seen” set to track visited node.
- While add each neighbor into stack, remove current node from neighbor’s neighbor list!

LC Graph#8: 323. Number of Connected Components in an Undirected Graph.

```
1 class Solution {
2     vector<vector<int>> G;
3     vector<bool> visited;
4 public:
5     void dfs(int i) {
6         visited[i] = true;
7         for (int j = 0; j < G.size(); j++) {
8             if (G[i][j] && visited[j] == false) dfs(j);
9         }
10    }
11
12    int countComponents(int n, vector<vector<int>>& edges) {
13        visited = vector<bool>(n, false);
14        G = vector<vector<int>>(n, vector<int>(n, 0));
15        for (int i = 0; i < edges.size(); i++) {
16            G[edges[i][0]][edges[i][1]] = 1;
17            G[edges[i][1]][edges[i][0]] = 1;
18        }
19
20        int res = 0;
21        for (int i = 0; i < n; i++) {
22            if (visited[i] == false) {
23                dfs(i);
24                res++;
25            }
26        }
27
28        return res;
29    }
30};
```



- Transform edges into G
- Use vector<bool> visited(n, false) to track if node is visited by dfs
- Start dfs with node 0, every return of dfs, increment res++;

LC Graph#9:

743: network delay time.

```

1+ class Solution {
2+ public:
3+     int networkDelayTime(vector<vector<int>>& times, int n, int k) {
4+         // establish G and adj list for each node
5+         int N = n + 1;
6+         vector<vector<int>> G(N, vector<int>(N,0));
7+         vector<vector<int>> adj(N); // holds neighbors for easy access
8+         for (int i = 0; i < times.size(); i++) {
9+             int u = times[i][0];
10+            int v = times[i][1];
11+            int w = times[i][2];
12+            G[u][v] = w;
13+            vector<int> neis1 = adj[u];
14+            neis1.push_back(v);
15+            adj[u] = neis1;
16+        }
17+
18+        vector<int> time(n+1, INT_MAX) ; // hold the final result
19+        vector<bool> visited(n+1, false);
20+        visited[0] = true; // since we are not using 0 position, exclude it
21+        time[k] = 0; // set the given node K as starting point
22+

```

44
45
46
47
48
49
50
51

You are given a network of n nodes, labeled from 1 to n . You are also given times , a list of travel times as directed edges $\text{times}[i] = (u_i, v_i, w_i)$, where u_i is the source node, v_i is the target node, and w_i is the time it takes for a signal to travel from source to target.

We will send a signal from a given node k . Return the time it takes for all the n nodes to receive the signal. If it is impossible for all the n nodes to receive the signal, return -1 .

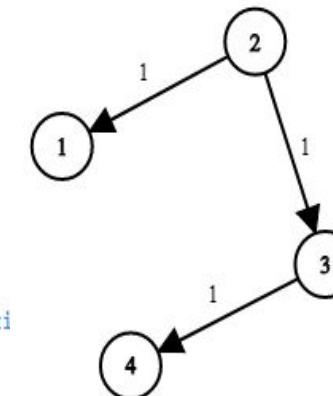
```

while (1) {
    // find the node u with mini time from not visited nodes
    // make u visited
    int u = 0;
    int minTime = INT_MAX;
    for (int i = 1; i < N; i++) {
        if (visited[i] == false & time[i] < minTime) {
            u = i; minTime = time[i];
        }
    }
    if (minTime == INT_MAX) break; // all nodes are covered
    visited[u] = true;
    // find all neighbors of u, named it v, and update time based on given time
    vector<int> u_neis = adj[u];
    for (int j = 0; j < u_neis.size(); j++) {
        int v = u_neis[j];
        if (visited[v] == false) {
            int w = G[u][v];
            time[v] = min(time[u] + w, time[v]);
        }
    }
}
int result = 0;
for (int i = 1; i < time.size(); i++)
    result = max(result, time[i]);

return result == INT_MAX? -1:result;
}

```

Example 1:



Input: $\text{times} = [[2,1,1],[2,3,1],[3,4,1]]$, $n = 4$, $k = 2$
Output: 2

visited T F T F F F F F

time 0 INF INF

k

- Among NOT visited nodes, find node with mini time as node u
- Find all neighbors of u , name as v , $\min(\text{time}[v], \text{time}[u]+w)$, make node u visited.
- Repeat above until all nodes are visited or TIME=INF(not reachable)

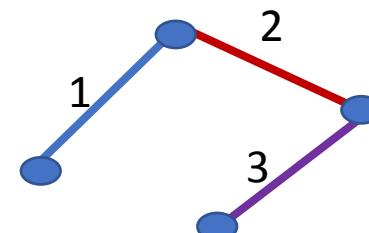
LC Graph#10: 787: cheapest flights within K stops.

- Key points

LC# 1135. Connecting Cities With Minimum Cost (Prims algorithm)

```
1 class Solution {
2 public:
3     // use prims algorithm:
4     // select the smallest edge, but make sure one node is in node hashset, the other one is not
5     // until N nodes in hashset
6     static bool cmp(vector<int> c1, vector<int> c2) {
7         return c1[2] < c2[2];
8     }
9
10    int minimumCost(int n, vector<vector<int>>& connections) {
11        int res = 0;
12        unordered_set<int> nodes;
13        sort(connections.begin(), connections.end(), cmp); // nLogN
14        // take the first minimum cost edge
15        vector<int> first = connections[0];
16        res = first[2];
17        nodes.insert(first[0]);
18        nodes.insert(first[1]);
19        connections.erase(connections.begin());
20        while (nodes.size() < n) {
21            bool foundConnection = false;
22            for (int i = 0; i < connections.size(); i++) { // N*N
23                vector<int> c = connections[i];
24                bool node1InSet = (nodes.find(c[0]) != nodes.end());
25                bool node2InSet = (nodes.find(c[1]) != nodes.end());
26                if ((node1InSet && !node2InSet) || (!node1InSet && node2InSet)) {
27                    res += c[2];
28                    if (!node1InSet) nodes.insert(c[0]);
29                    if (!node2InSet) nodes.insert(c[1]);
30                    connections.erase(connections.begin() + i);
31                    foundConnection = true;
32                    break;
33                }
34            }
35            if (!foundConnection && (nodes.size() < n)) return -1;
36        }
37        return res;
38    }
39}
```

- Sort connection first
- Add first mini cost connection, remove it from connections, add nodes into set
- Find mini cost connection with only one node in set until all nodes are in.



Misc: LRU cache, LC#146

Tree Height, Prime number and Inorder Traversal

```
int height_dfs(Node* node, int &res) {
    if (node == NULL) return -1;
    if (node->children.size() == 0) return 0;

    vector<int> heights;
    for (auto n: node->children) {
        heights.push_back(height_dfs(n, res));
    }
    sort(heights.begin(), heights.end(), greater<int>());
    if (heights.size() >= 2) {
        res = max(res, heights[0] + heights[1] + 2); // add two: one for each side
    }
    return heights[0]+1;
}

class Solution {
public:
    int countPrimes(int n) {
        if (n <= 2) return 0;

        vector<bool> prime_nums(n, true);
        for (int p = 2; p*p < n; p++) {
            if (prime_nums[p]) {
                for (int i=p*p; i < n; i +=p)
                    prime_nums[i] = false;
            }
        }
        int res = 0;
        for (int i = 2; i < n; i++)
            if (prime_nums[i]) res++;

        return res;
    }
};
```

```
vector<int> inorderTraversal(TreeNode* root) {
    vector<int> res;
    if (root == NULL) return res;
#define RECURSIVE
    vector<int> left = preorderTraversal(root->left);
    for (int i = 0; i < left.size(); i++)
        res.push_back(left[i]);
    res.push_back(root->val);
    vector<int> right = preorderTraversal(root->right);
    for (int i = 0; i < right.size(); i++)
        res.push_back(right[i]);
#define ENDIF
    TreeNode *cur = root;
    stack<TreeNode *> st;
    while (cur != NULL || !st.empty()) {
        if (cur != NULL) {
            st.push(cur);
            cur = cur->left; // push down all the way to bottom most left node
        } else {
            cur = st.top(); st.pop();
            res.push_back(cur->val);
            cur = cur->right;
        }
    }
    return res;
}
```

From Excel

- 937s, 408s, 65h, 242s, 49m, 1062m, 1092m, 14m, 12m, 527h
- 752m(open lock, BFS), 56s, 80m, 88s, 238s, 16m, 41m, 128m, 35s,
- 69s, 34s, 153m, 162m(peak), 278m, 33m, 81m, 34m, 50m, 96m,
- 762h, 136s, 137m, 169s, 229s, 134m(gas), 179m, 402m, 55m, 45m,
- 31m, 19s, 21s, 82s, 83s, 86s, 206s, 92s, 61m, 109m,
- 138m, 141s, 142m, 143m, 148m, 538s, 110s(balanced tree), 102s, 107s(Tree Level T), 144s (Tree Preorder T),
- 104s, 105m(binT from Pre & inroder), 297m, 98m, 285m, 510m, 366m, 156m(BT upDown), 39m, 207m(course schedule),
- 51m, 52m, 90m, 78m, 47m, 46m, 63s, 70s, 53s, 152m(max product subarray),
- 72m, 115h, 120m(trangle min sum), 139m(word break, recur with memo), 140h (word break II, backtrack), 97m, 91m, 128m(int array LongConSeq), 23m,
- 232m, 155m, 263m, 264m(ugly number), 212h, 79h, 295h, 84h, 438s, 311m, 288m,
- 981m(time based key store), 706m(Design HashMap)

Java: HashMap and HashSet

```
HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();  
/* Add new entry */  
map.put(1, 100); map.put(2, 200); if (map.size() == 2) System.out.println("Cool");  
  
/* Get/check using key */  
Integer val = map.get(1);  
/* Loop thru key */ for (Integer key: map.keySet()) System.out.println("key is " + key.intValue());  
/* Loop thru value */  
for (Integer val: map.values()) System.out.println("val is " + val.intValue());  
  
/* Iterator using Map.Entry*/  
for (Map.Entry<Integer, Integer> ent: map.entrySet())  
    System.out.println("Key is "+ ent.getKey() + ", value is " + ent.getValue());  
  
HashSet<String> set = new HashSet<String>();  
set.add("xyz"); if (set.contains("abc")) System.out.println("abc is not found");
```

Java: Stack and Queue

```
Stack<Integer> st = new Stack<Integer>;
st.push(1); st.push(2); st.pop(); st.peek(); st.size();
while (!st.isEmpty()) System.out.println("pop: " + st.pop());
```



```
Queue<Integer> q = new LinkedList<Integer>();
q.add(1); q.add(2); q.add(3);
System.out.println("Q size is " + q.size());
while (!q.isEmpty()) System.out.println("Q remove : " + q.remove()); // remove() throw exp if empty
Integer val = q.poll(); //remove() method returns the head of the queue and removes it.
                    // It returns null if the queue is empty.
```



```
Vector<Integer> v = new Vector<Integer>(); // thread safe
v.add(1); v.add(2); Integer val = v.get(1);
v.remove(0); v.remove(new Integer(2));
for (int i = 0; i < v.size(); i++) System.out.print(v.get(i) + " ");
```



```
ArrayList<Integer> al = new ArrayList<Integer>(); // not synchronized, faster
al.add(1); al.add(2); Integer val = al.get(1);
al.remove(0); al.remove(new Integer(2));
for (int i = 0; i < al.size(); i++) System.out.print(al.get(i) + " ");
```

Java: util.Arrays

- `Java.util.Arrays.binarySearch(int[] arr, int key):`
- `Java.util.Arrays.copyOf(int[] arr, int newLen)`
- `Java.util.Arrays.sort(int[] arr) or`
`Java.util.Arrays.sort(int[] arr, Collections.reverseOrder())`
- Class MyCmp implements Comparator<Student> {
 public int compare(Student a, Student b) { return a.id – b.id; } }
- `Java.util.Arrays.sort(Student[] arr, new MyCmp())`

Python: list (just like vector in C++ std)

- L1 = [1,2,3] L2 = list()
- L1.append(4) L2.append("abc")
- L1.insert(0, 100) # add 100 at the beginning (index, val)
- val1 = L1.pop(0) # remove first element and return
- Val2 = L1.pop() # pop last element, for this reason, list can be used as stack
- L1.remove(2) # remove the first occurrence of value 2
- L1.reverse()
- L1.sort()
- L1.count(2) // tell how many occurrence of value 2
- L1.index(3) □ L1[3]// return value at given index
- Size = len(L1) # return how many elements in the list

Python dict: (unordered_map in C++ std)

- D1 = { 1:"abc", 2:"xyz"} D2 = dict()
- D1[100] = “odfag” D2[“xyz”] = 987 // add new element
- **len(D1)** # get the total elements in dictionary
- **del D1[2]** #remove element using key
- D1.pop(2) # another way to remove element using key
- 2 **in** D1: # tell if key exists
- “abc” **in** D1.**values()**: # tell if value exists
- for k in D1: print(k); print(D1[k]) # iteratre thru all elements

Python queue & stack

```
from queue import Queue  
q = Queue(maxsize = 3)  
q.put('a'); q.put('b'); q.put('c')  
print(q.get())  
q.empty()  
q.full()  
q.qsize()
```

```
from queue import LifoQueue  
st = LifoQueue(maxsize = 3)  
st.put('a'); st.put('b'); st.put('c')  
print(st.get())  
st.empty()  
st.full()  
st.qsize()
```

