Artificial Intelligence 4 Games

Pathfinding Overview

Pathfinding

Basic Formulation

- → Given a graph, a start node, a goal node
- → Our task is to find the shortest (cheapest) path from the start to the goal
- → In more general case, edges have weights
- → Many more sophisticated variants

Applications

All types of movement:

★ Units in strategy games, NPCs in role playing games, all entities in simulation/tycoon games, enemies in FPS,

. . .

Basic Algorithms

- ★ Breadth First Search
- ★ Dijkstra
- ★ Best-First Search
- **★** A*

Uninformed Search

Breadth First Search (BFS)

- → We expand equally in all directions
- → The edges need to have uniform weights
- → Worst case performance O(|V|+|E|)
- → Complete (always finds the path) and optimal (found path is the best one)

Dijkstra's algorithm

- → We expand lower cost paths first
- → Works with weighted edges
- ⇒ Worst case performance $\Theta(|V|\log|V|+|E|)$
- → Complete and optimal

Informed Search

Best-First (Greedy) Search

- → We use heuristic estimation of how far is from a given node to the goal. E.g. on a plane the straight line distance
- → Expand nodes that "should be" closest to the goal
- Mostly get to the goal fast using some nonoptimal path
- → In the worst case it can explore the whole graph like a badly guided DFS
- → Requires cycle checking to be complete

A*

- \rightarrow We combine uniform-cost approach ("backward cost" g(x)) with the best-first approach given by the heuristic h(x)
- Thus, the A* algorithm use the ordering based on f(x) = g(x) + h(x).
- → Works similar as Dijkstra(if h(x)=0 then it simply is Dijkstra)
- → Complete
- → Optimal with a proper heuristic



Pseudocode

- closedSet := {}
- openSet := {START}
- g(x) = 0 if x==START, otherwise ∞
- f(x) = h(x) if x==START, otherwise ∞
- while openSet is not empty:
 - current := openSet.dequeue_lowest_f()
 - if current==GOAL: return Success.
 - closedSet.add(current)
 - for each neighbor of current:
 - if neighbor in closedSet: continue
 - varf := g(current)+ cost(current, neighbor) + h(neighbor)
 - **if** neighbor **not in** openSet **or** f(neighbor) > varf:
 - openSet.add(neighbor)
 - f(neighbor) = varf
 - return Failure

Heuristic

The heuristic function has to be admissible (optimistic), i.e. $h(x) \leftarrow true_cost(x, goal)$.

To ensure that A* finds optimal solution, the heuristic function has to be consistent (monotone), i.e. f is not decreasing along any path: $h(x) \le true_cost(x, x') + h(x')$

All consistent heuristics are admissible.

Pathfinding in Games

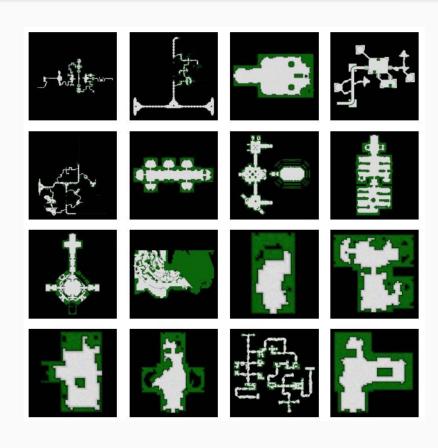
Literature

- ★ Botea, A., Bouzy, B., Buro, M., Bauckhage, C., Nau, D. <u>Pathfinding in games</u>. Dagstuhl Follow-Ups, vol 5, pp. 21-31, 2013.
- ★ Abd Algfoor, Z., Sunar, M. S., Kolivand, H. <u>A</u>
 comprehensive study on pathfinding techniques for robotics and video games. International Journal of Computer Games Technology, 2015.
- ★ Sturtevant, N., GPPC: Grid-Based Path Planning Competition: http://movingai.com/GPPC/

Pathfinding in Games

- hierarchical planning,
- → non-trivial heuristic functions,
- → dynamic changes in the environment,
- → multiple targets,
- → multi-agent pathfinding,
- → adversarial pathfinding,
- various types of terrains and mobilities of the units,
- → incomplete information,
- → real-time constraints, memory constraints,
- → inventory-driven pathfinding,
- → various graph types (grid, hex, navmesh),
- interpretation: doorways are nodes / doorways are edges,

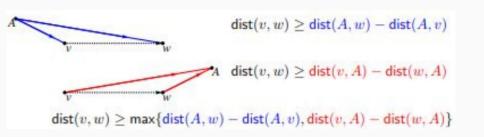




Landmark Heuristics

ALT

- → We choose a small set of landmark points
- → And precompute distances between each node and landmark
- → We can use those distances as a heuristic, as the maximum distance difference over a subset of landmarks is admissible



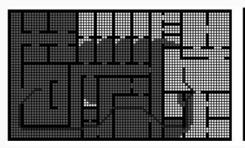
ALTBest_P

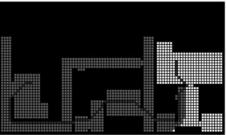
- → We have a predefined set of P landmarks,
- → During a search we choose a single landmark that gives the highest h value for the root node (initially, the start point)
- → The heuristic is maximum of the ALT value for the selected point and the Manhattan heuristic
- → ALTBest_p is worse in terms of quality than ALT, but it is cheaper per node.

Room-based Heuristics

Dead-end heuristic

- → In preprocessing phase, we decompose map into disjoint areas
- → During a search query, we start with identifying and removing from consideration irrelevant areas (by setting the heuristic values to ∞)

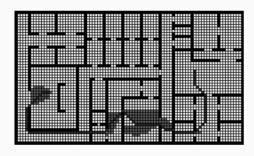




Gateway heuristic

- → We decompose map into areas, borders between areas form gates
- → Then we precompute distances between all gates
- → And use gate distance within a heuristic

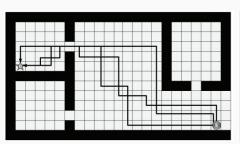
$$h^G(n,g) = \min \sum_i \sum_j h^l(n,G_i) + H(G_i,G_j) + h^l(G_j,g)$$

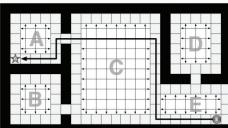


Symmetry Elimination

One of the main problems with A* pathfinding on grid maps is related to open spaces, where exist exponential number of optimal paths.

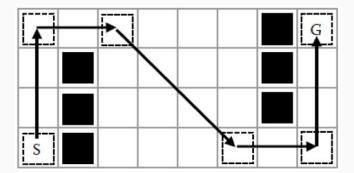
This can be solved by e.g. decomposing maps into obstacle-free rooms and pruning all nodes except the ones on the perimeter.





JPS

- → Ultrafast A* improvement over the uniform cost octile grids
- → Uses successor-pruning technique to remove redundant paths from consideration
- → Replace neighbors with further away nodes and jumps directly to them omitting opening intermediate nodes



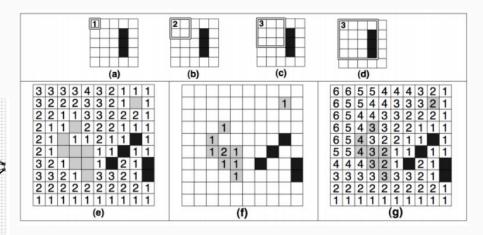
Hierarchical Pathfinding

HPA*

- → In preprocessing step we decompose map into disjoint square sectors
- → We put some entrances on the edges, and compute true distances between entrances within one sector
- → Pathfinding uses A* on the abstract graph,
- → Refinement step to smooth the path

HAA*

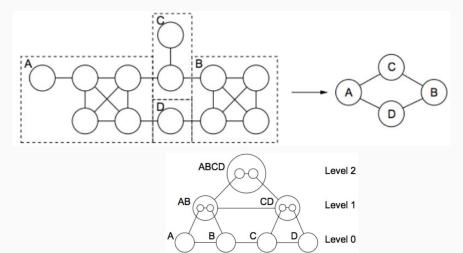
- → Extension of HPA*
- → Allow several types of terrain
- → Handles units of variable sizes and different terrain traversal capabilities



Hierarchical Pathfinding

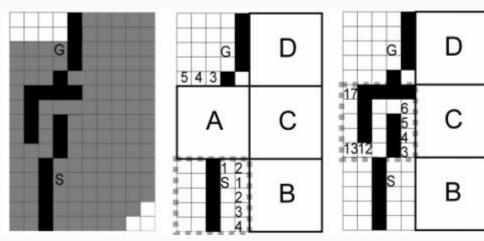
PRA*

- → Builds bottom-up hierarchical representation by merging small fully connected regions
- → We search top-down, common parent mean path between regions



Block A*

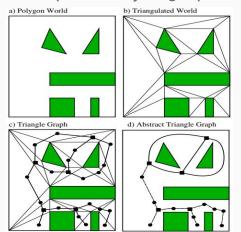
- → We precompute a database of all mxn block topologies
- → And for each block compute distances between all boundary nodes

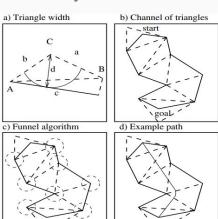


Triangulation-based Environments

TA*

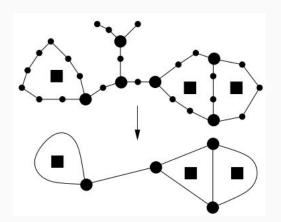
- → Maps are represented with obstacles defined by polygons
- → Uses A* like algorithm on graphs induced by the triangulation
- → Finds optimal any-angle paths for circular objects







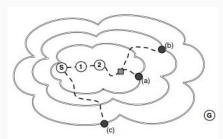
- → TA* improvement reducing the triangulation graph to contain only nodes of degree 3
- → Runs much faster



Real-time Search

TBA*

- → Standard A* that can be interrupted
- → The most promising node on the open list is traced back to the start
- → If it passes through the agent's position he simply follows the path
- Otherwise, the agent backtracks his steps towards the start
- → Or use shortcut-search enhancement

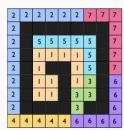


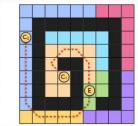
AA*, ARA*

- → A* is called multiple times, with h value multiplied by gradually decreasing ε≥1
- → Repair variant reuses and updates the path

LRTA*, D LRTA*

- → Planning, learning, acting steps
- → Partitioning on regions, computed paths between
- → Entries to regions set as subgoals



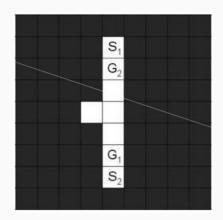




Cooperative Pathfinding

Cooperative A*

- → Agents reserves their paths
- → They can wait until the path is free.
- → Reservation table is hashable
- → The initial ordering of the agents has huge impact on the outcome.



Push and Swap

- → Suboptimal but complete
- → Introduces two high-level operations.
- → Push: forcing all other agents to get out the way

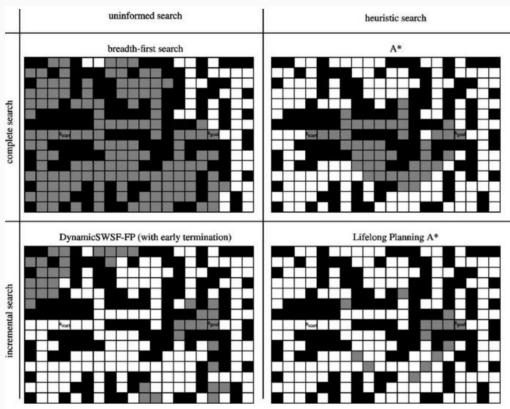
Dynamic Environment

LPA*

- → Incremental version of A*.
- Tracks changes between distance from the start of the node's predecessors
- → Forces recalculations when required

D*, Focused D*, D*Lite

- → Search backwards (from goal to start)
- → Node can be: new/open/closed/raise/lower



Playing Games with A*

Literature

- ★ Tristan Penman: N-puzzle https://tristanpenman.com/demos/n-puzzle/
- ★ André G. Pereira, Marcus Ritt, Luciana S. Buriol: Optimal Sokoban solving using pattern databases with specific domain knowledge, Artificial Intelligence, vol. 227, pp. 52-70, 2015.
- ★ Julian Togelius, Sergey Karakovskiy, Robin Baumgarten:

 The 2009 Mario Al Competition. IEEE CEC, 2010.

A* as a Game Playing Algorithm

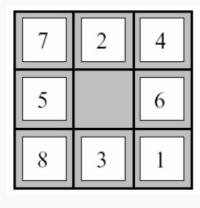
A* can be also used for playing (i.e. solving a game). Then our "map" is actually a graph of game states, and "moves" are legal actions.

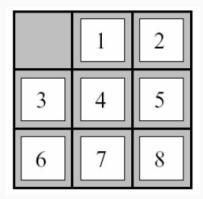
This is particularly popular approach when dealing with single-player games, which turns "playing" into a planning problem.

A few examples of games that can be dealt with A*:

- → N-puzzle
- → Sokoban
- → Mario
- → Codingame Fall Challenge 2020 :-)

8-Puzzle





Heuristics

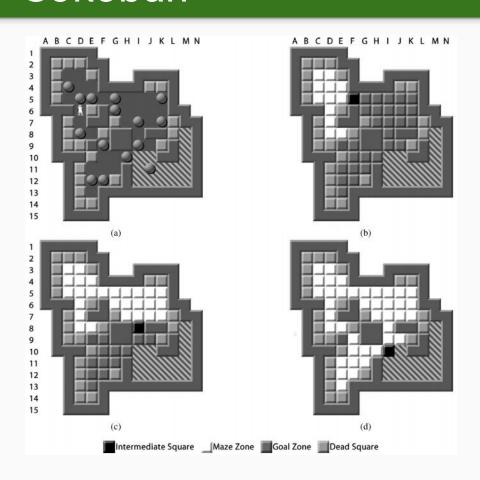
Popular method of generating heuristics is relaxation of the problem.

We used that with pathfinding: straight-line distance is simply no-obstacle assumption.

- ★ What If we could just swap any two tiles?
 Number of misplaced tiles.
- ★ What if we could slide any direction at any time? *Total Manhattan distance*.

We usually got trade-off between the quality of the heuristic (less nodes expanded) and its computational cost.

Sokoban



- → IDA*
- → Quite complicated heuristic
- → Instance decompositions
- → Pattern databases
- → Dead squares detection
- → Domain-dependent tie-breaking rules
- → ...

Mario



- → Goal is to reach the right side of the screen
- → As fast as possible
- → Interruptible A*: best node so far is used
- → Slight heuristic overestimation
- → Recalculating plan every two game ticks

<u>Video</u>

Bonus material: Video games

StarCraft

- http://www.codeofhonor.com/blog/the-starcraft-path-finding-hack
- ★ https://youtu.be/I9mCau4a130
- ★ https://youtu.be/paX8nHGPpXA

Command & Conquer

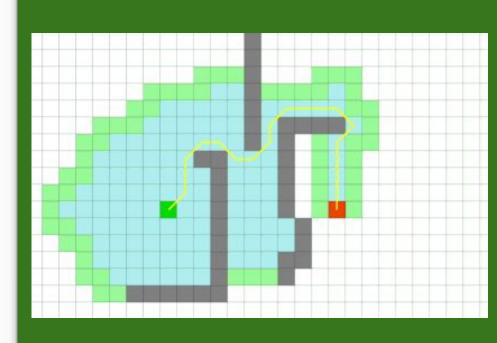
★ https://youtu.be/Wb84Vi7XFRg?t=525

Cities: Skylines / Sim City

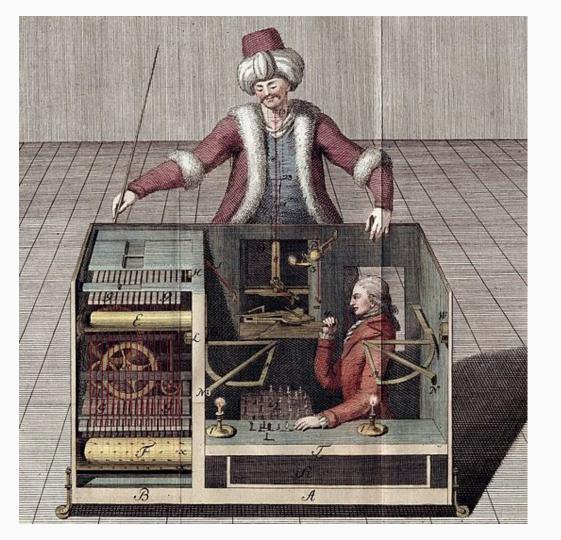
- ★ https://youtu.be/MeNxJVOL9eM
- ★ https://youtu.be/zHdyzx_ecbQ

Summary

- ★ Pathfinding is actually a complex task
- ★ (We just scratched the surface of the problem)
- ★ It is also one of the most important tasks
- ★ Most algorithms are based on A*



Thanks!



Bonus reference quiz

