

# Introduction to Computer Networks



## IEEE 802.3 Ethernet

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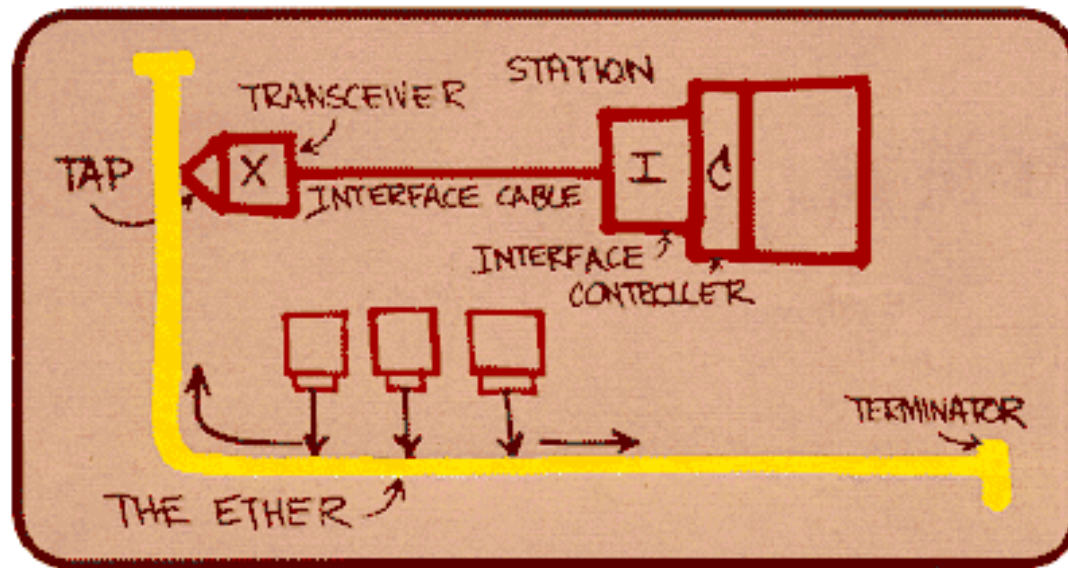
# Outline

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- **Introduction**
- **Ethernet Topologies**
- **Ethernet Frame Format**
- **Ethernet MAC Protocol -- CSMA/CD**
- **802.3 Ethernet Standards**

# Ethernet

- Most successful local area networking technology of last 30 years.
- First widely used LAN technology
- kept up with speed race: 10 Mbps – 100 Gbps



Metcalfe's Ethernet sketch

# Ethernet

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- Developed in the mid-1970s by researchers at the Xerox Palo Alto Research Centers (PARC).
- DEC and Intel joined Xerox to define a **10-Mbps Ethernet** standard in 1978.
- This standard formed the basis for **IEEE standard 802.3**
- More recently 802.3 has been extended to include
  - 100-Mbps version called **Fast Ethernet**,
  - 1000-Mbps version called **Gigabit Ethernet**,
  - **10 Gigabit Ethernet**, and also
  - 100 Gigabit Ethernet

# Ethernet: Unreliable, Connectionless

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- **Connectionless:** No handshaking between sending and receiving NICs
- **Unreliable:** receiving NIC doesn't send ACKs or NACKs to sending NIC
- Ethernet's MAC protocol: **Carrier Sense Multiple Access with Collision Detection (CSMA/CD)**

不确认对方正确接收

# Outline

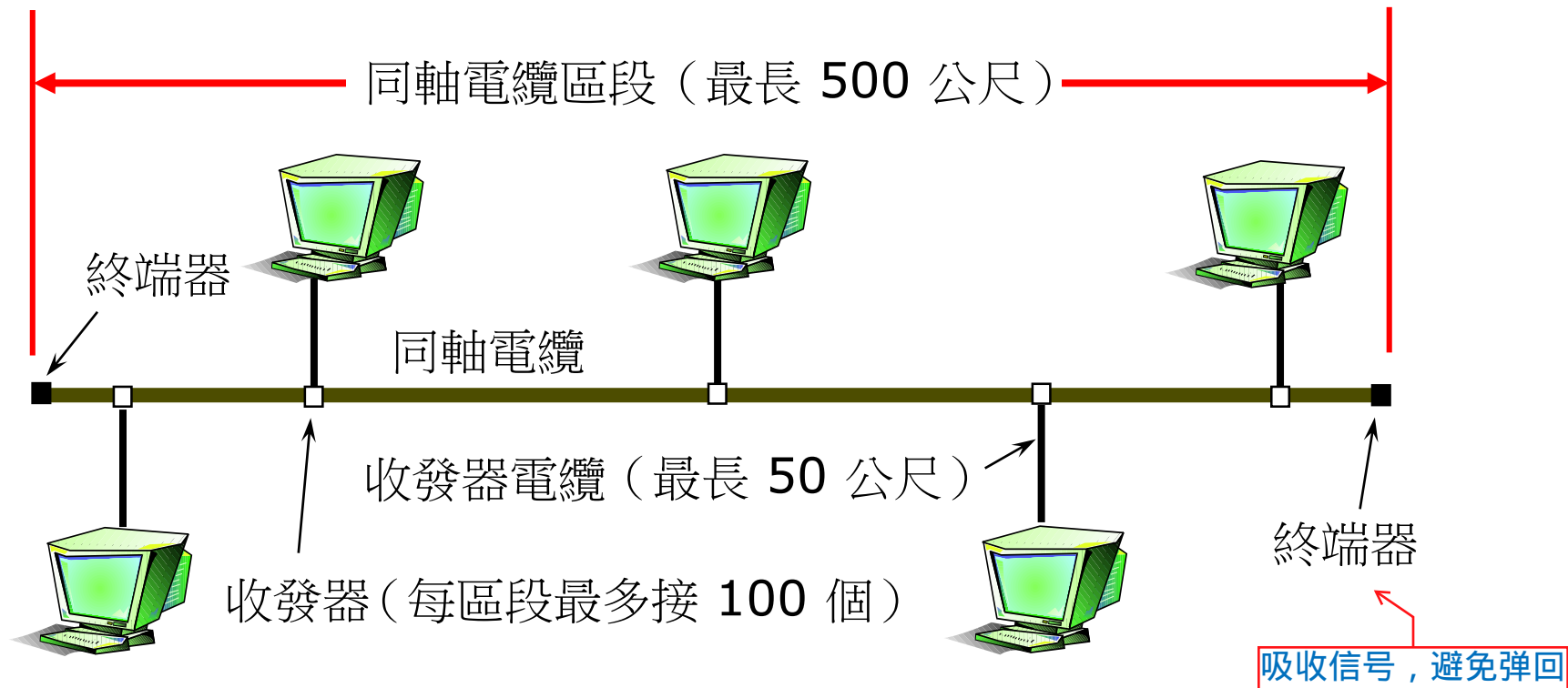
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# Bus Topology

## ■ Bus topology popular through mid 90s

- all nodes in same collision domain (can collide with each other)

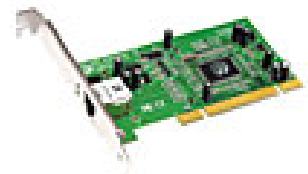


# Ethernet (10Base5)

网段, 最长500m

- An Ethernet segment is implemented on a coaxial cable of up to **500 m**.
- Hosts connect to an Ethernet segment by tapping into it.
- A **transceiver** (a small device directly attached to the tap) detects when the line is idle and drives signal when the host is transmitting.
- The transceiver also receives incoming signal.
- The transceiver is connected to an **Ethernet adaptor** which is plugged into the host. But now most are built in into the computers.
- The protocol is implemented on the adaptor.

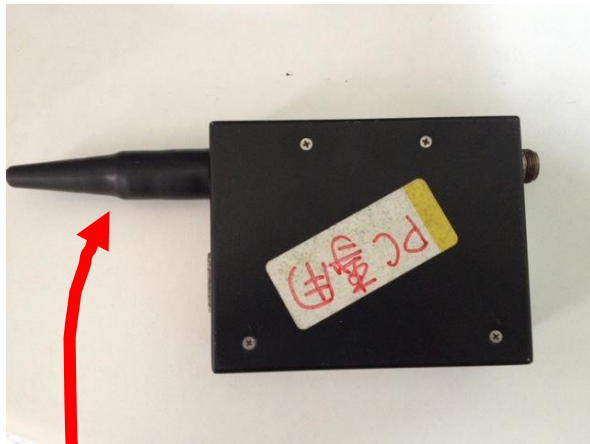
既可传送也可接收



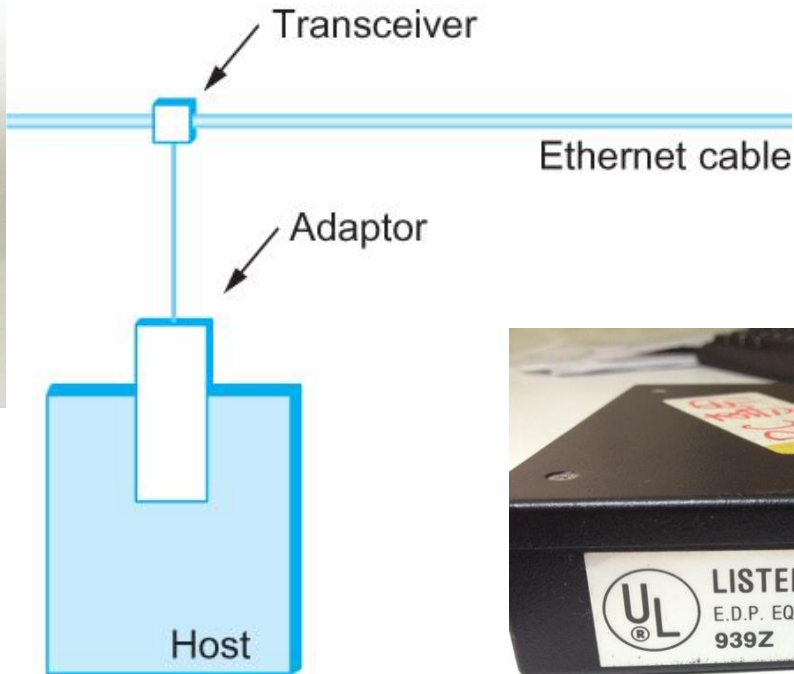


# Ethernet (10Base5)

10Mbps, Base技术只有一个channel, 最长500m

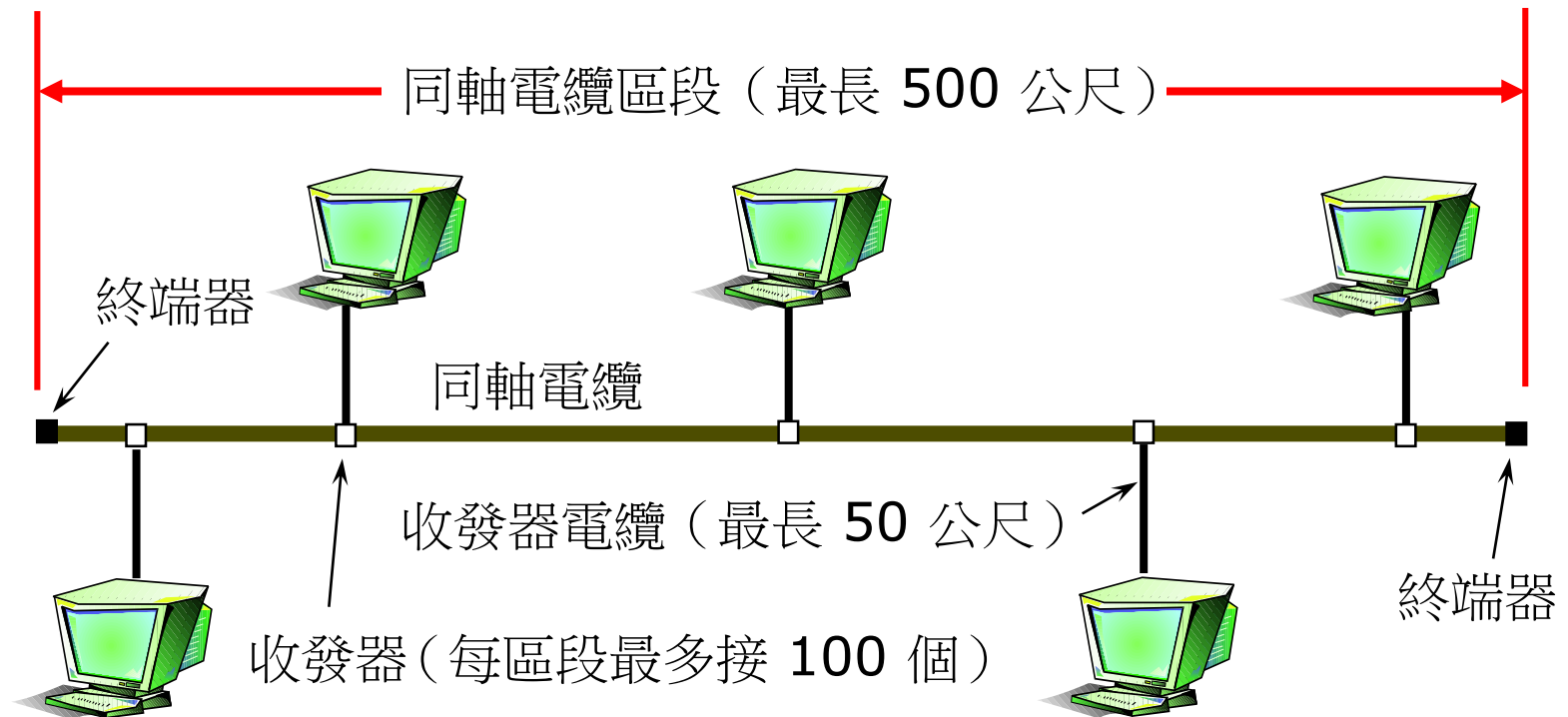


Terminator

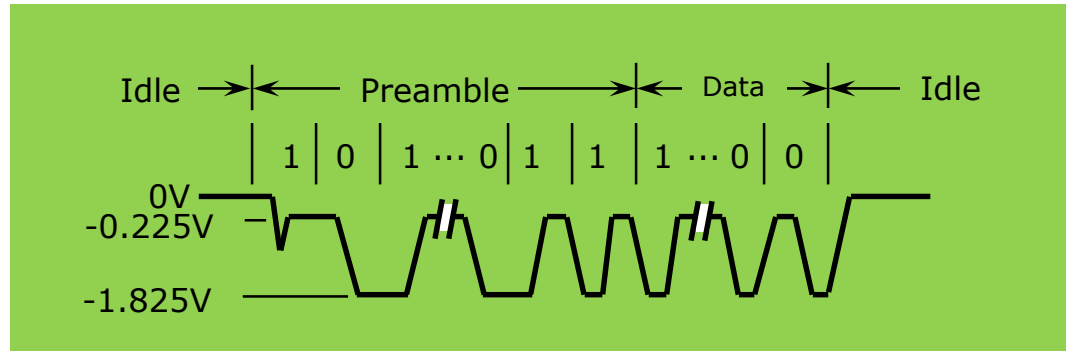


Ethernet transceiver, adaptor, and terminator

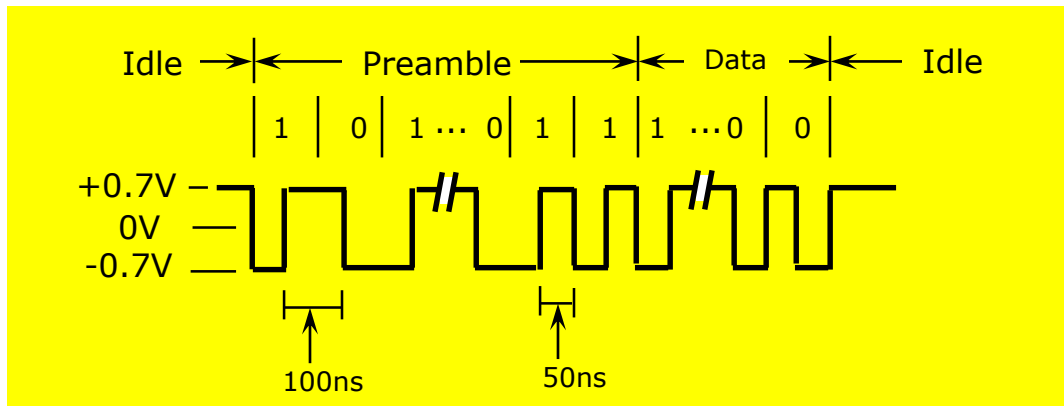
# Network Configuration Example 1 (Single segment)



# Cable Signaling (Manchester Encoding)



Coaxial Cable



Transceiver Cable

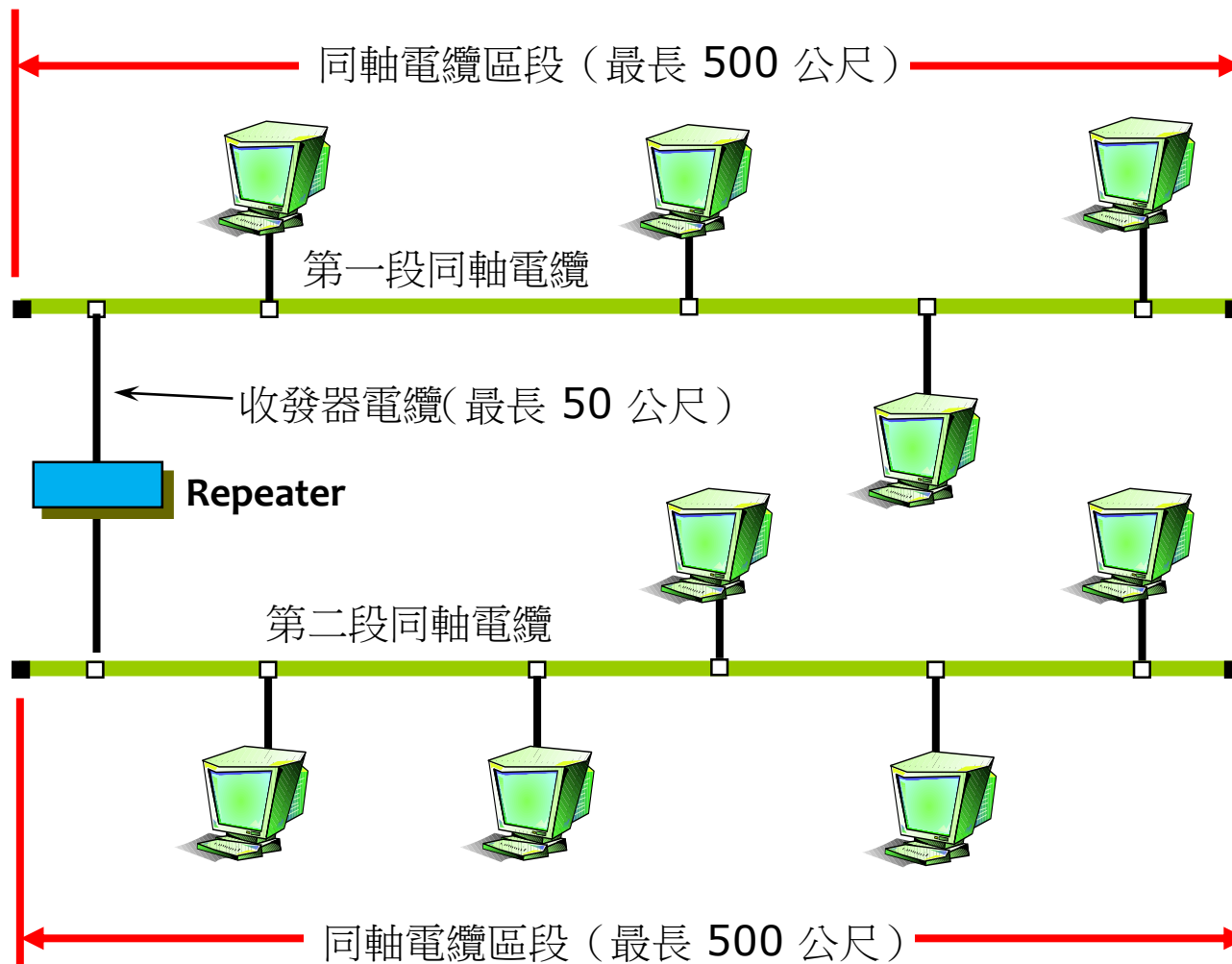
- Each bit has a transition
- Allows clocks in sending and receiving nodes to synchronize to each other

# Ethernet (10Base5)

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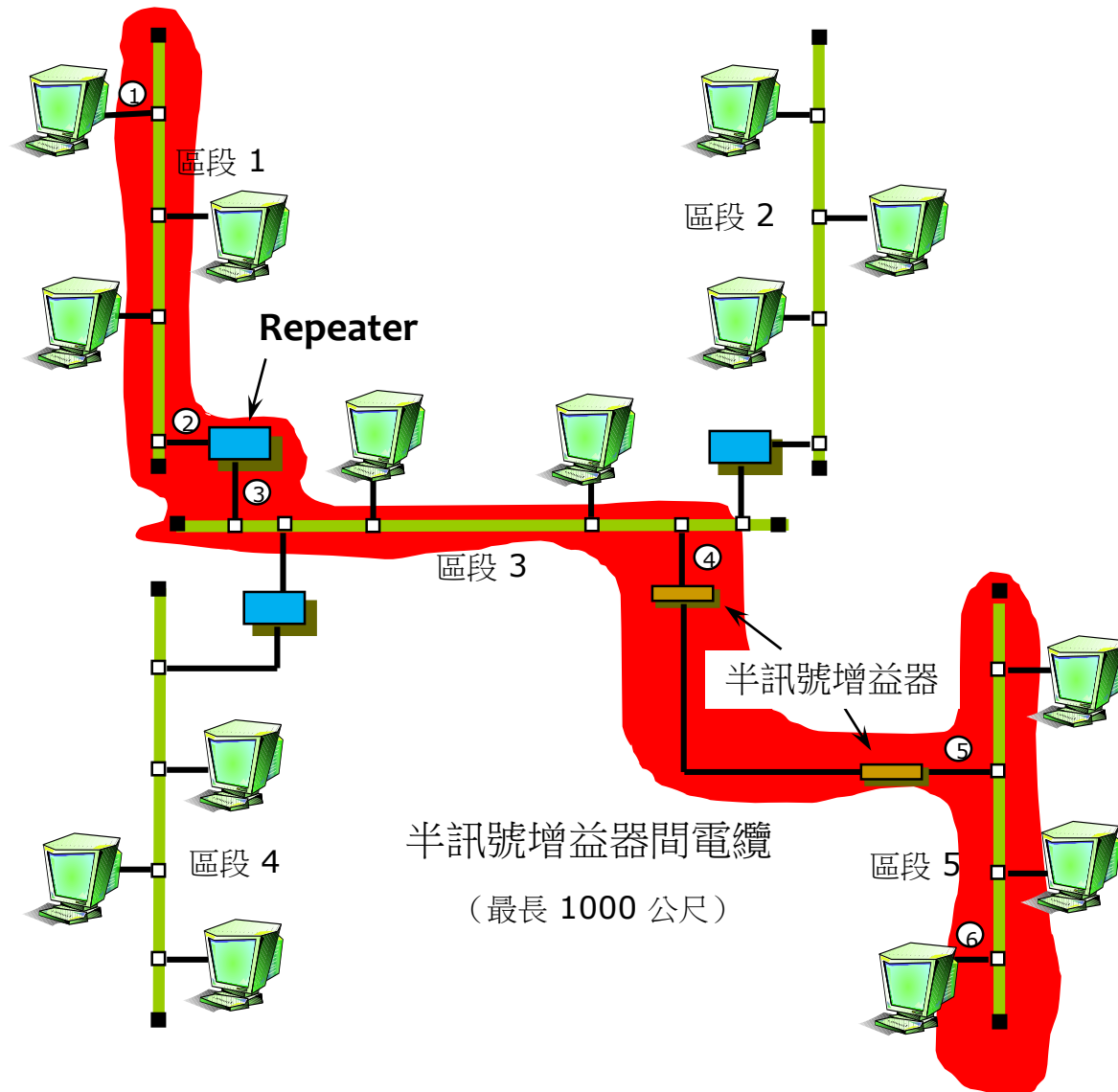
- Multiple Ethernet segments can be joined together by **repeaters**.  
 第一层设备, 只负责转送讯号而不管MAC地址
- A **repeater** is a device that forwards **digital signals**.  
 将衰减的信号恢复成方波
- No more than **four repeaters** may be positioned between any pair of hosts.
  - An Ethernet has a total reach of only 2500 m.

# Network Configuration Example 2 (Two segments)



# Network Configuration Example 3

## (Five segments, maximum)



# Ethernet (10Base2)

## ■ New Technologies in Ethernet

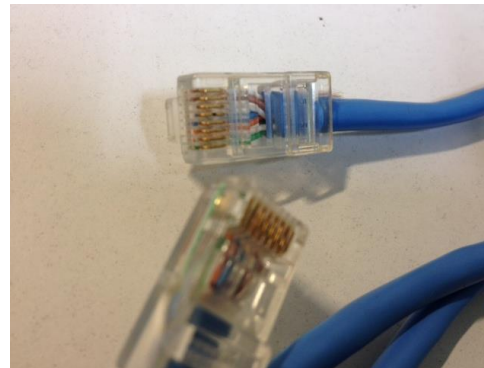
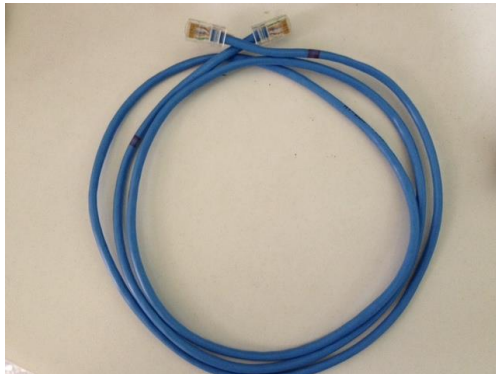
- Instead of using coax cable, an Ethernet can be constructed from **a thinner cable** known as **10Base2** (the original was 10Base5) ↗ 更轻便，传输距离也更短
  - ▶ **10** means the network operates at 10 Mbps
  - ▶ **Base** means the cable is used in a baseband system
  - ▶ **2** means that a given segment can be no longer than 200 m



# Ethernet (10BaseT)

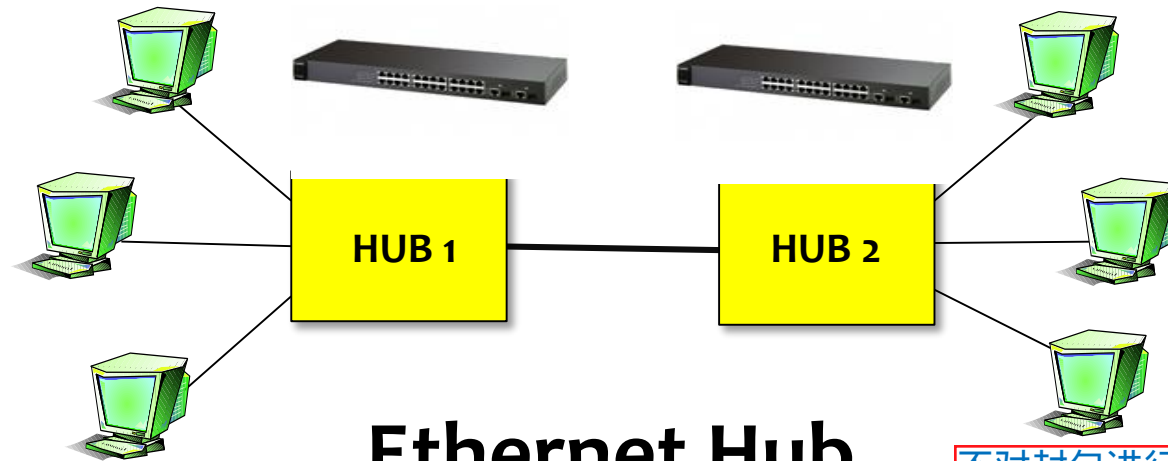
## ■ New Technologies in Ethernet

- Another cable technology is 10BaseT
  - ▶ T stands for twisted pair 双绞线
  - ▶ Limited to 100 m in length
- With 10BaseT, the common configuration is to have several **point to point segments** coming out of a multiway repeater, called **Hub** 管理收纳



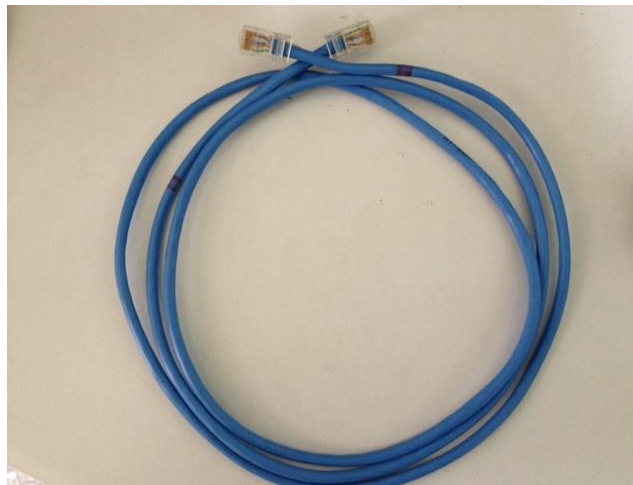


# Ethernet



**Ethernet Hub**

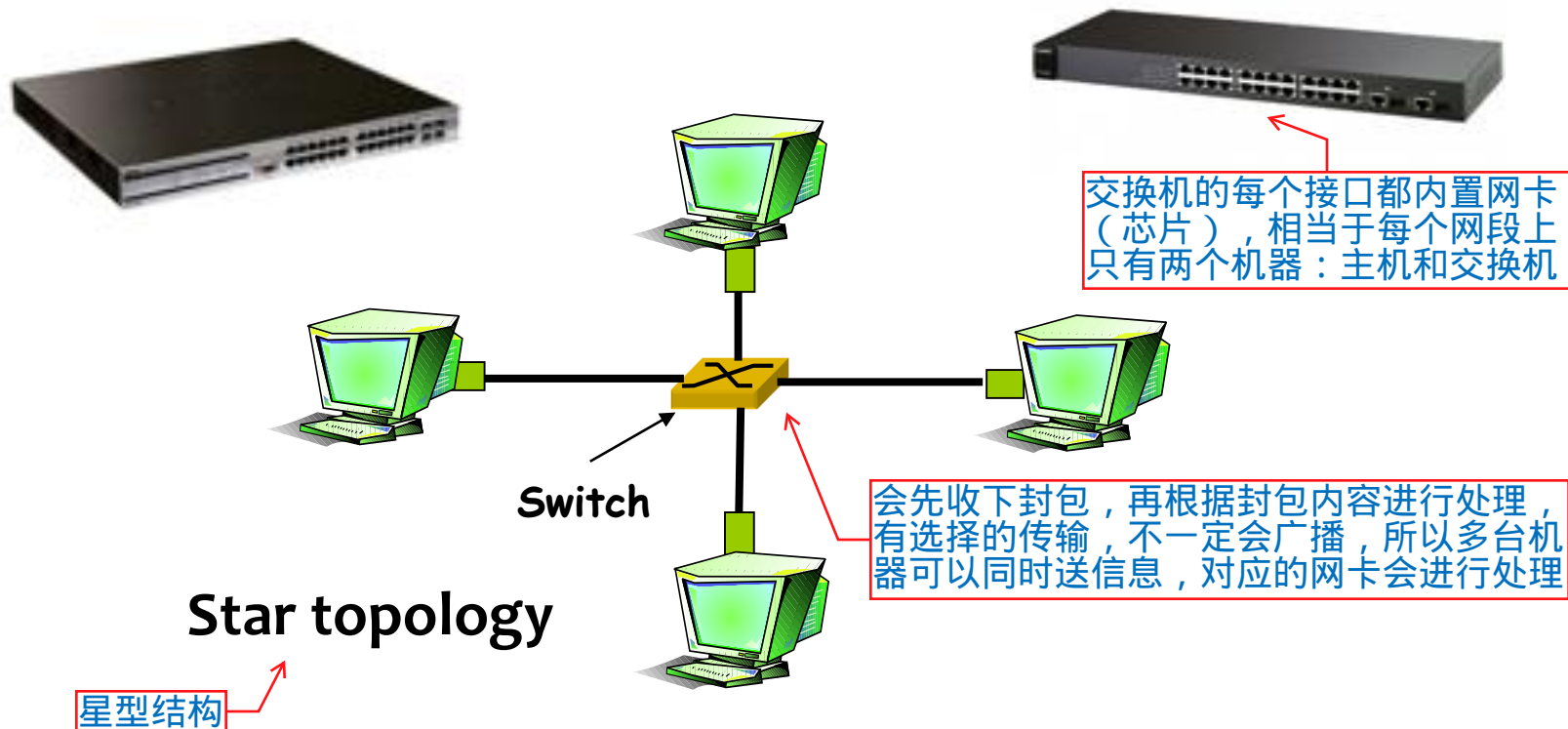
不对封包进行处理，收到后即广播出去，相当于Repeater，是第一层设备。如果多个机器同时送会造成冲突



# Star Topology

## ■ Today: **Star topology** prevails

- active **switch** in center
- each “spoke” runs a (separate) Ethernet protocol (nodes do not collide with each other)



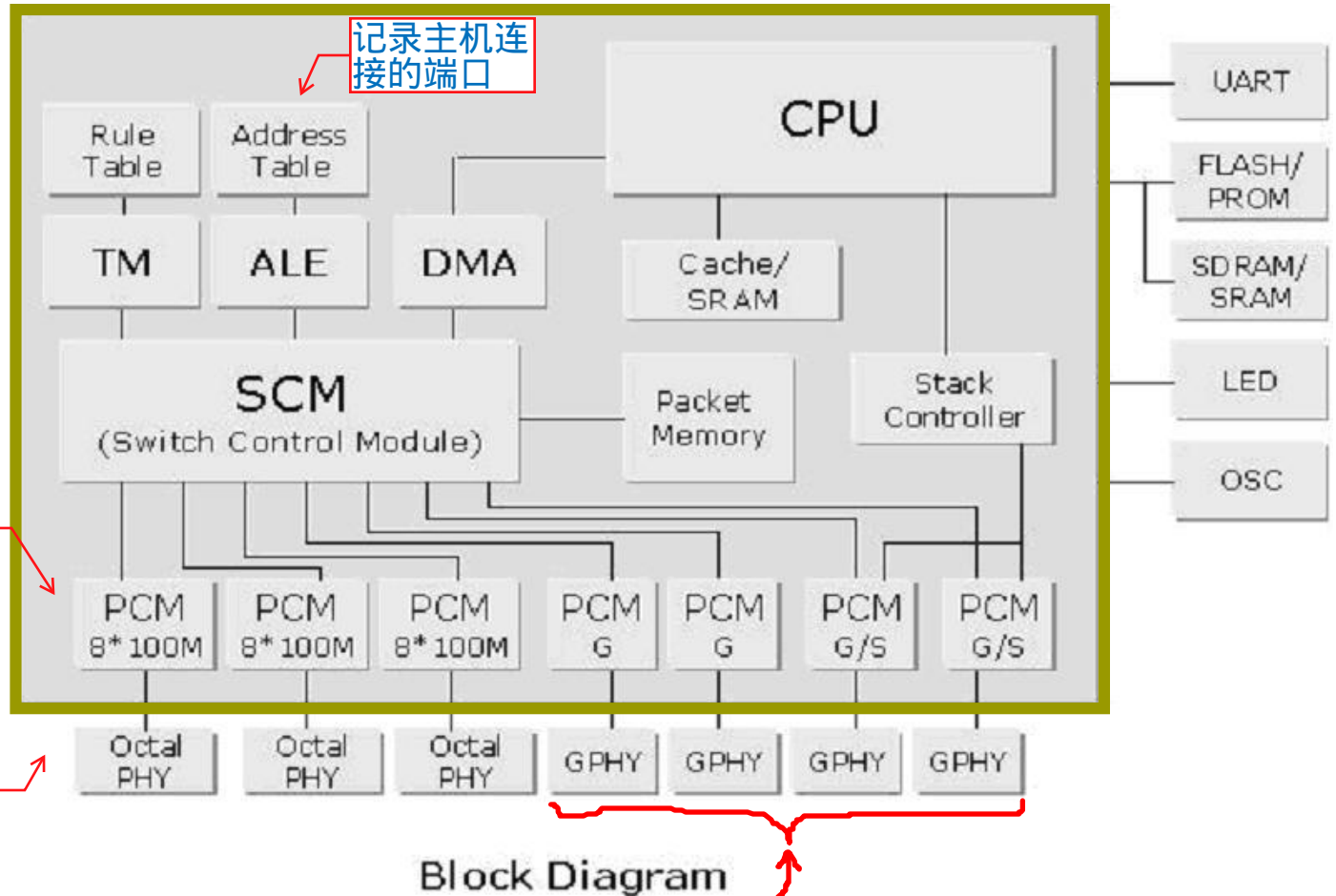
# Ethernet-Switch

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- To speed up the transmission rate of Ethernet Hub without changing the interface cards on stations.
- Ether-Switch Architecture
- Each Ethernet port can have a transmission simultaneously.



# Ethernet Switch ASIC example



## Acute Leo AQ6628 24+4 Ethernet Switch ASIC

每个封包进来，同时查表然后同时交换出去->速度很快而且所有工作都在芯片完成；能够以封包进来的速度完成处理并转送出去->while speed 线速

# Outline

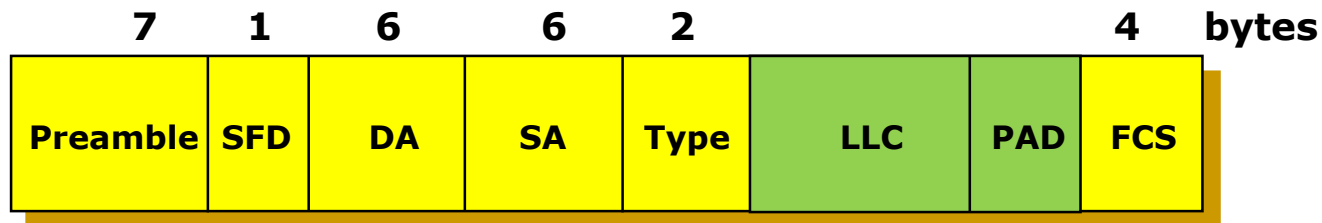
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- **Ethernet Frame Format**
- Ethernet MAC Protocol -- CSMA/CD
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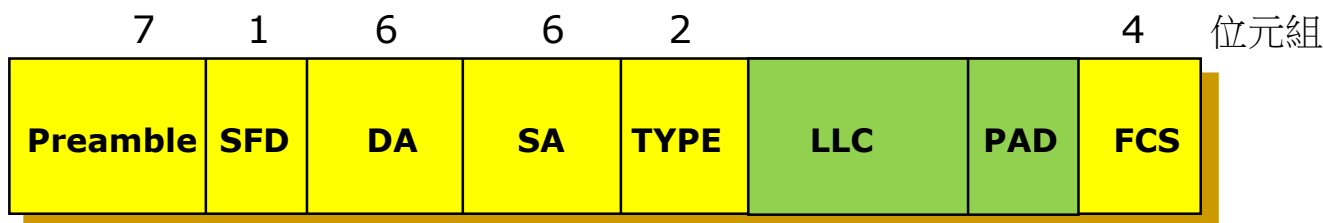
# Ethernet Frame Format

## ■ Frame format

- Preamble (64bits): allows the receiver to synchronize with the signal (sequence of alternating 0s and 1s). 网卡的ID具有全球唯一性
- Source and Destination MAC Addresses (48bits each).
- Packet type (16bits): acts as demux key to identify the higher level protocol. 出于公平
- Data (up to 1500 bytes)
  - ▶ Minimally a frame must contain at least 46 bytes of data.
  - ▶ Frame must be long enough to detect collision.
- FCS: CRC (32bit)



# Ethernet Frame Format



- Preamble: (101010...1010) for Synchronization
- SFD: Start Frame Delimiter (10101011)
- DA: Destination MAC Address
- SA: Source MAC Address
- Packet type (16bits): acts as demux key to identify the higher level protocol.
- LLC-Frame: Up to 1500 bytes
- PAD: Padding when LLC-Frame < 46 bytes
- FCS: Frame Check Sequence (CRC-32)
- MAC-frame size -- from DA to FCS

Cycl i cal Redundancy Check 循环冗余检验

内容至少46byte+首尾18byte

- **Min 64 bytes to distinguish from collision**
- **Max 1518 bytes to prevent dominating bandwidth**

内容最多1500byte，首尾18byte不包括Preamble&SFD

# Ethernet Addresses

- Each host on an Ethernet (in fact, every Ethernet host in the world) has a **unique Ethernet Address**.
- The address belongs to the adaptor, not the host.
  - It is usually burnt into ROM.
- Ethernet addresses are typically printed in a human readable format
  - As a sequence of six numbers separated by colons.
  - Each number corresponds to 1 byte of the **6 byte address** and is given by a pair of hexadecimal digits, one for each of the 4-bit nibbles in the byte
  - Leading 0s are dropped.
  - For example, **8:0:2b:e4:b1:2** is

地址是属于网卡的，而非主机，  
不插网卡的主机是没有地址的

十六进制

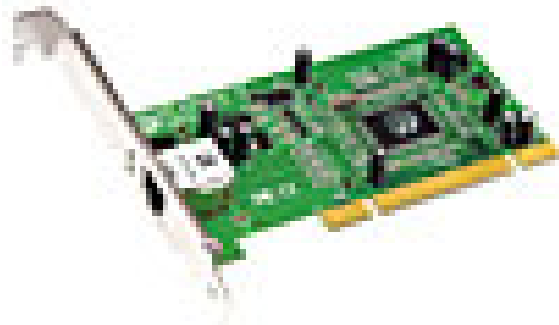
▶ 00001000 00000000 00101011 11100100 10110001 00000010



# Ethernet Addresses

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- To ensure that every adaptor gets a unique address, each manufacturer of Ethernet devices is allocated a **different prefix** that must be prepended to the address on every adaptor they build
  - ▶ AMD has been assigned the 24bit prefix **8:0:20**



# Ethernet Addresses

广播

- Each frame transmitted on an Ethernet is received by **every adaptor** connected to that Ethernet.
- Each adaptor recognizes those frames addressed to its address and passes only those frames on to the host.
- In addition to **unicast address**, an Ethernet address consisting of all 1s is treated as a **broadcast address**.
  - All adaptors pass frames addressed to the *broadcast* address up to the host.
- Similarly, an address that has the first bit set to 1 but is not the *broadcast* address is called a **multicast address**.
  - A given host can program its adaptor to accept some set of *multicast* addresses.

DA的48位全为1

# Ethernet Addresses

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- To summarize, an Ethernet adaptor **receives all frames** and accepts
  - Frames addressed to its own address
  - Frames addressed to the broadcast address
  - Frames addressed to a multicast address if it has been instructed

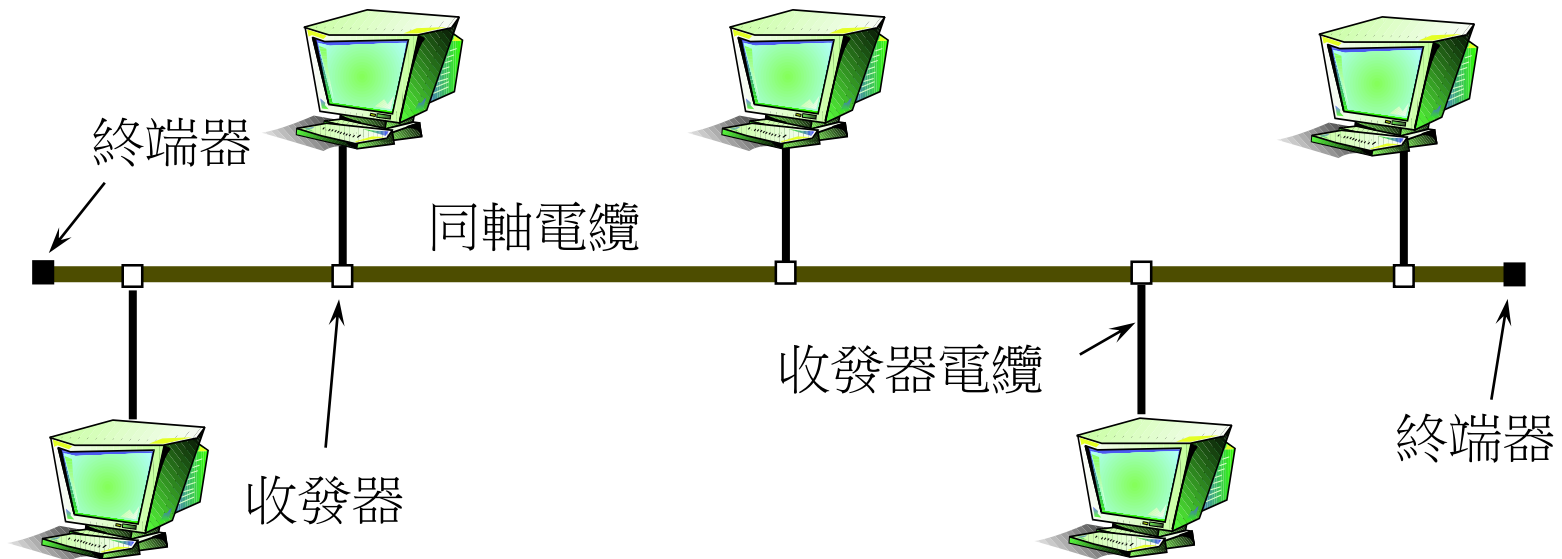
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# Ethernet MAC protocol

- Any signal placed on the Ethernet by a host is **broadcast over the entire network**
  - Signal is propagated in both directions.
  - Repeaters forward the signal on all outgoing segments.
  - **Terminators** attached to the end of each segment absorb the signal.



# CSMA (Carrier Sense Multiple Access)

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**CSMA:** listen before transmit:

If channel sensed idle: transmit entire frame

■ If channel sensed busy, defer transmission

# CSMA collisions

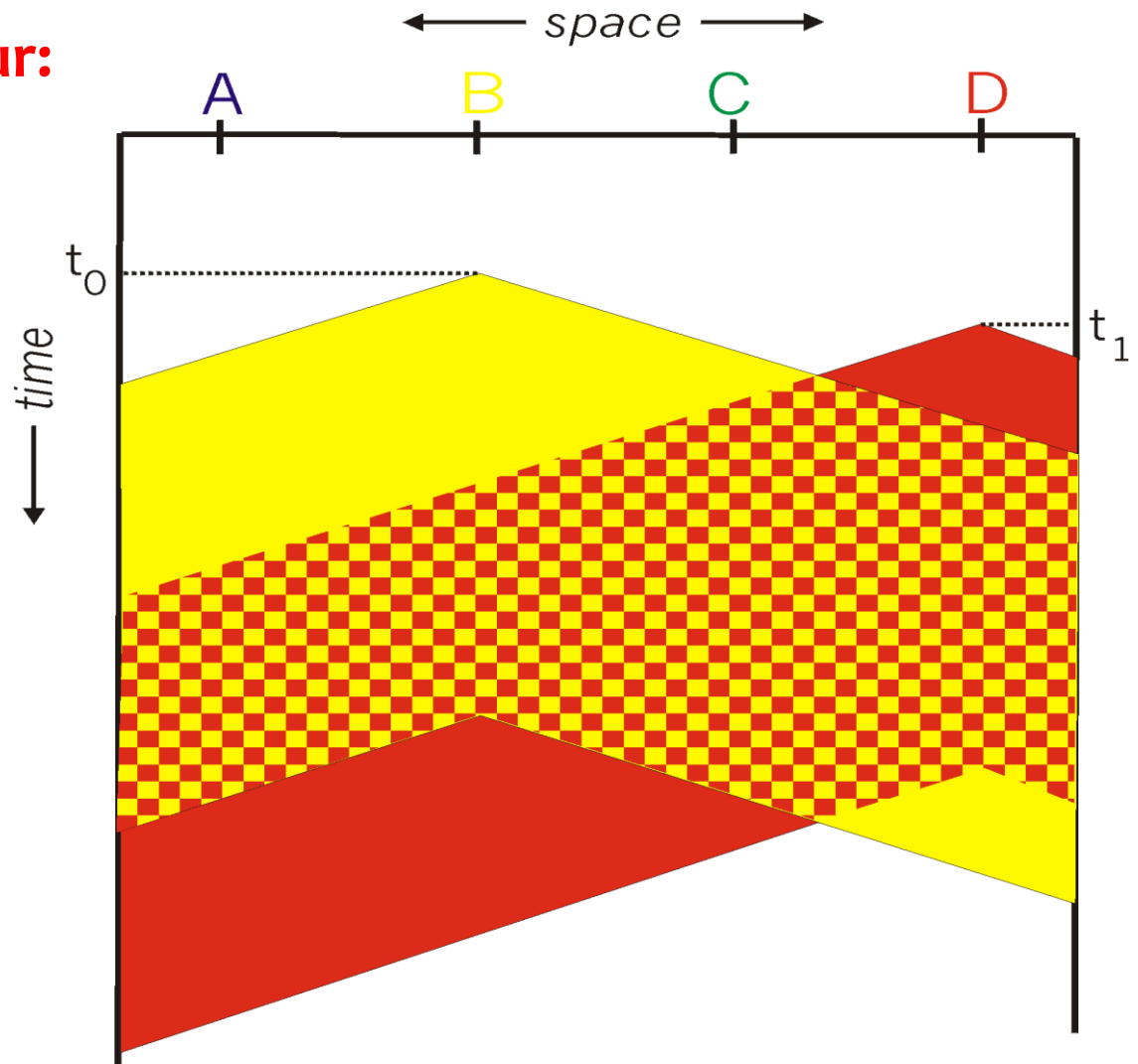
**collisions can still occur:**

propagation delay means  
two nodes may not hear  
each other's transmission

**collision:**

entire packet transmission  
time wasted

发送以后继续听，发现冲撞  
不作反应仍然将封包送完



# CSMA/CD (Collision Detection)

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**CSMA/CD:** carrier sensing, deferral as in CSMA

- collisions *detected* within short time
- **colliding transmissions aborted**, reducing channel wastage

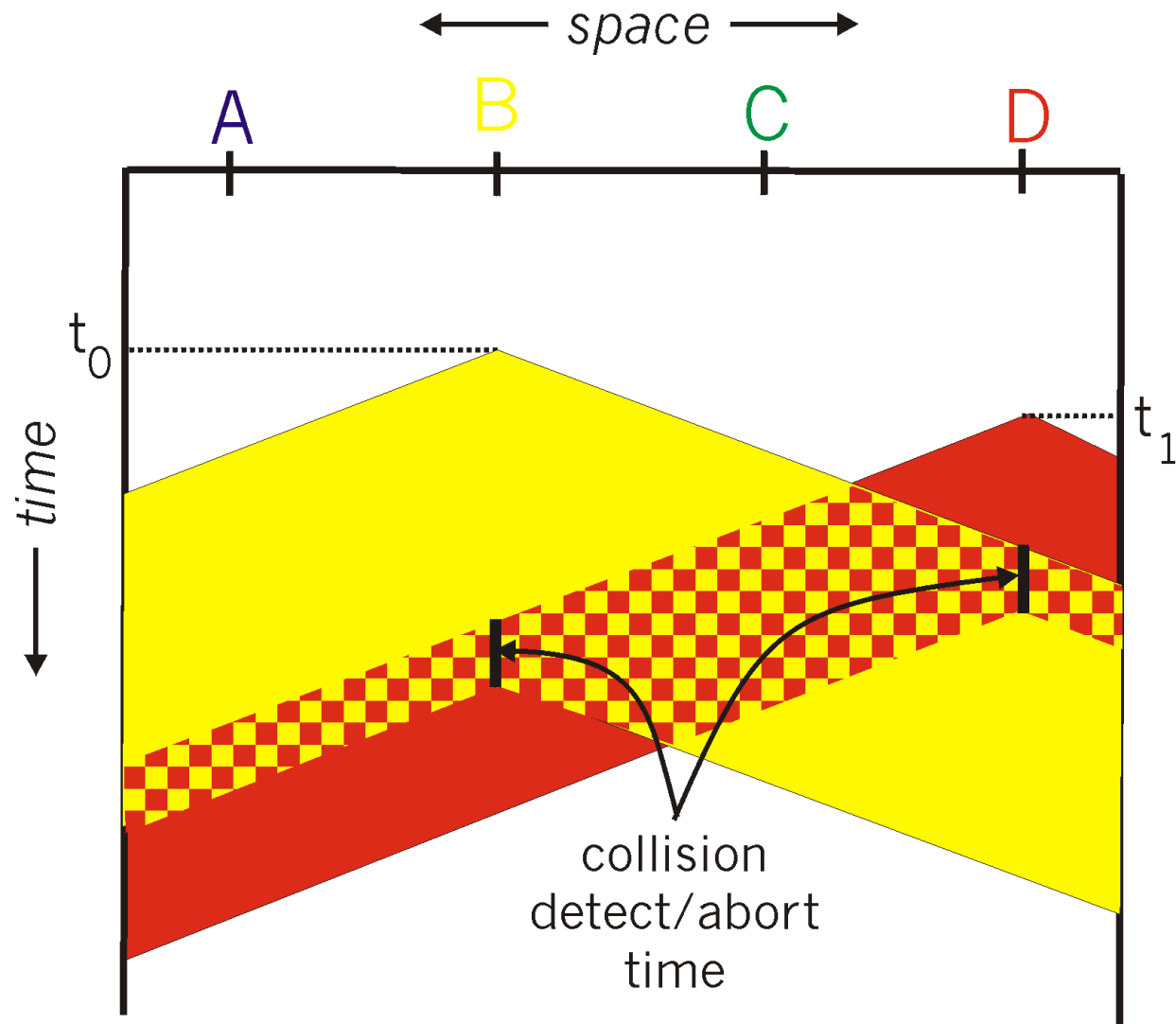
↑ 检测到封包冲撞即停止传送

## ■ Collision detection:

- Measure signal strengths, compare transmitted, received signals



# CSMA/CD collision detection



# CSMA/CD

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- Carrier Sense Multiple Access with Collision Detection (**CSMA/CD**).
  - A set of nodes send and receive frames over a shared link.
  - **Carrier sense** means that all nodes can distinguish between an **idle** and a **busy** link.
  - **Collision detection** means that a node listens as it transmits and can therefore detect when a frame it is transmitting has collided with a frame transmitted by another node.

# CSMA/CD

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- When the adaptor has a frame to send and the **line is idle**, it transmits the frame immediately.
- When the adaptor has a frame to send and the **line is busy**, it waits for the line to go idle and then transmits immediately.
- The Ethernet is said to be **1-persistent protocol** because an adaptor with a frame to send transmits with **probability 1** whenever a busy line goes idle.

监听到封包传送结束后发起传输的可能性

# CSMA/CD

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- Since there is no centralized control it is possible for two (or more) adaptors to begin transmitting at the same time,
  - Either because both found the line to be idle,
  - Or, both had been waiting for a busy line to become idle.
- When this happens, the two (or more) frames are said to be *collide* on the network.

# CSMA/CD

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- Since Ethernet supports **collision detection**, each sender is able to determine that a collision is in progress.
- At the moment an adaptor detects that its frame is colliding with another, it first makes sure to transmit a **32-bit jamming sequence** and then stops transmission.
  - Thus, a transmitter will minimally send 96 bits in the case of collision
    - ▶ 64-bit preamble + 32-bit jamming sequence

确保所有工作站都知道  
这个封包已经被冲撞

# CSMA/CD

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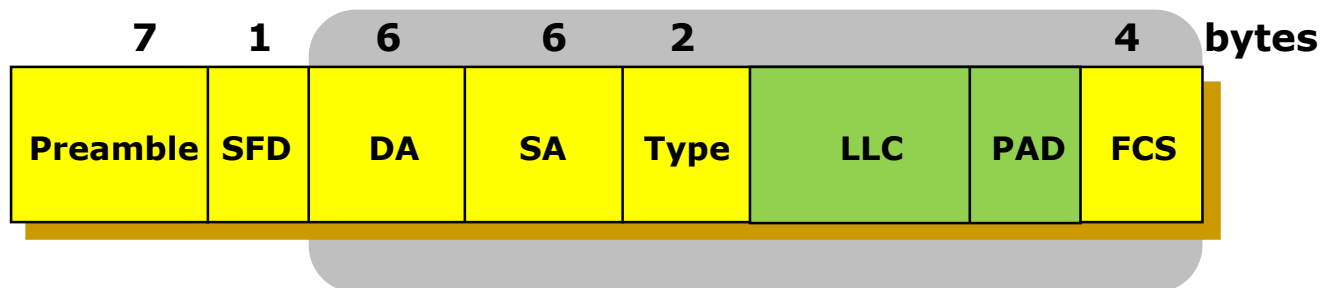
- One way that an adaptor will send only **96 bits** (called a ***runt frame***) is if the two hosts are **close to each other**.
- In case the two hosts are farther apart, they would have had to transmit longer, and thus send more bits, before detecting the collision.

# Collision Window

- The worst case scenario happens when the two hosts are **at opposite ends of the Ethernet**.
- To know for sure that the frame its just sent did not collide with another frame, the transmitter may need to **send as many as 512 bits**.
- Every Ethernet frame must be **at least 512 bits (64 bytes) long**.

若封包短于此长度，则不足够确认信号到网络另一端再回来是否发生了冲撞

- ▶ 14 bytes of header + 46 bytes of data + 4 bytes of CRC



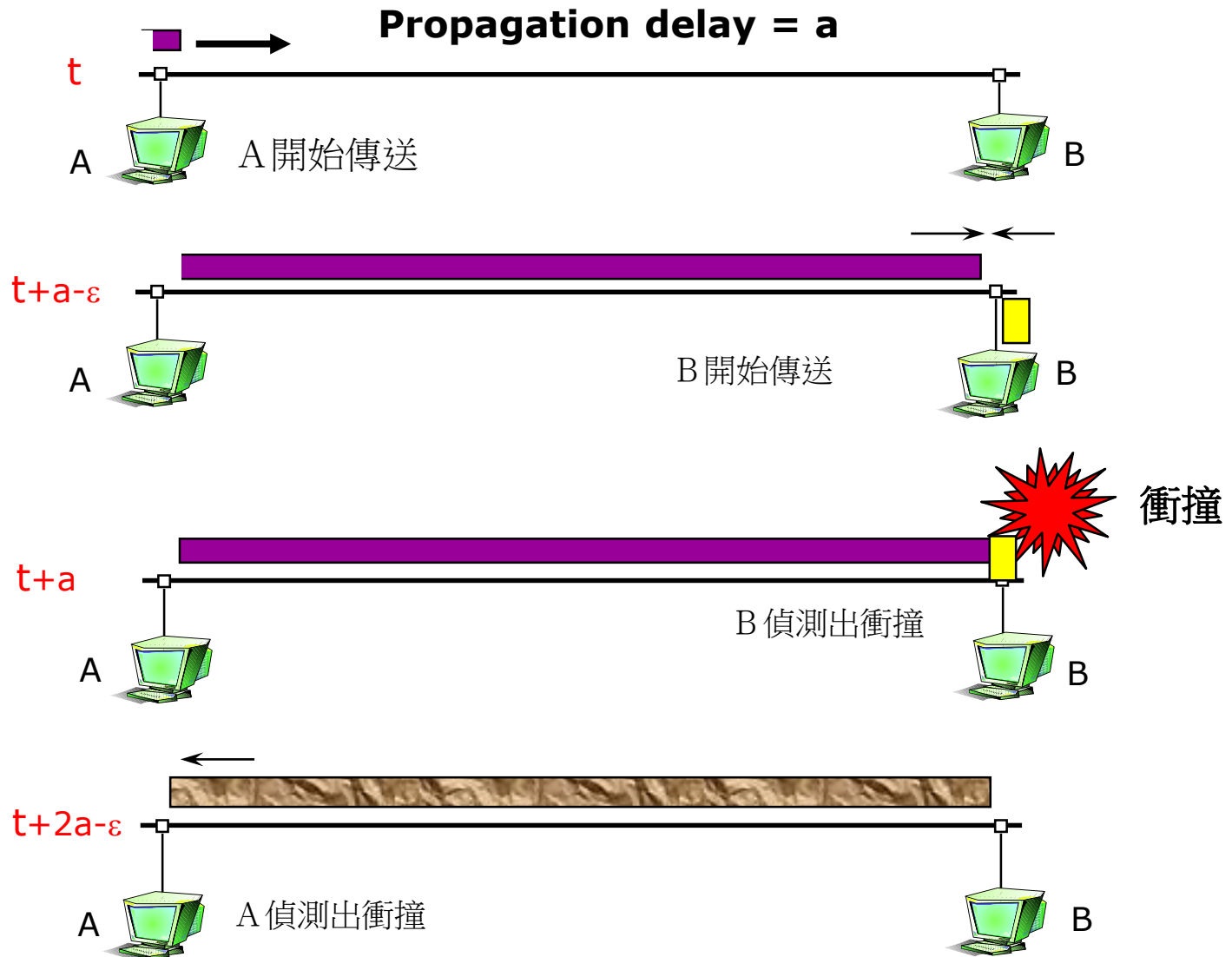
# Collision Window

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- **Why 512 bits (64 bytes) ?**
  - Why is its length limited to 2500 m?
- **Collision Window = round-trip delay ( $2a$ )**
- **The farther apart two nodes are, the longer it takes for a frame sent by one to reach the other, and the network is vulnerable to collision during this time**



# Collision Detection Window for CSMA/CD (=2a)



# Collision Window

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- A begins transmitting a frame at time  $t$
- $a$  denotes the one link latency
- The first bit of A's frame arrives at B at time  $t + a$
- Suppose an instant before host A's frame arrives, host B begins to transmit its own frame
- B's frame will immediately collide with A's frame and this collision will be detected by host B
- Host B will send the 32-bit jamming sequence
- Host A will not know that the collision occurred until B's frame reaches it, which will happen at  $t + 2a$
- Host A must continue to transmit until this time in order to detect the collision
  - Host **A must transmit for  $2a$**  to be sure that it detects all possible collisions

# Collision Window

- Consider that a maximally configured Ethernet is 2500 m long, and there may be up to four repeaters between any two hosts, the **round trip delay has been determined to be 51.2  $\mu$ s**
  - Which on **10 Mbps** Ethernet corresponds to 512 bits
  - **10 Mbps  $\times$  51.2  $\mu$ s = 512 bits**  $\leftarrow$  64byte
- The other way to look at this situation,
  - We need to limit the Ethernet's maximum latency to a fairly small value (51.2  $\mu$ s) for the access algorithm to work
    - ▶ Hence the maximum length for the Ethernet is on the order of 2500 m.

# Exponential Backoff Algorithm

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- Once an adaptor has detected a collision, and stopped its transmission, it **waits a certain amount of time and tries again**.
- Each time the adaptor tries to transmit but fails, it **doubles the amount of time it waits** before trying again.
- This strategy of doubling the delay interval between each retransmission attempt is known as ***Exponential Backoff***.

# Exponential Backoff Algorithm

- The adaptor first delays either 0 or 51.2  $\mu\text{s}$ , selected at random.
- If this effort fails, it then waits 0, 51.2, 102.4, 153.6  $\mu\text{s}$  (selected randomly) before trying again;
  - This is  $k * 51.2$  for  $k = 0, 1, 2, 3$
- After the third collision, it waits  $k * 51.2$  for  $k = 0 \dots 2^3 - 1$  (again selected at random).
- In general, the algorithm randomly selects a  $k$  between 0 and  $2^n - 1$  and waits for  $k * 51.2 \mu\text{s}$ , where  $n$  is the number of collisions experienced so far.

最多冲撞16次，16次都冲撞就放弃。但 $2^{16}$ 太大了，控制等待时间最长为 $2^{10}-1$ ，即冲撞10次以后就都在0-1023之间挑选等待时间

# CSMA/CD Protocol

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- Carrier Sense before transmission
- Carrier Sense while transmission
- Collision: Two or more stations transmitting simultaneously
- Backoff: Random delay after collision
- Deference: Defers transmission if channel is sensed busy
- **Collision Window (Slot time):** Round-trip propagation delay time plus some carrier sense time. In IEEE 802.3, this value is defined to be 51.2 us.

# CSMA/CD Collision Handling

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- Collision Signal is generated by Physical layer.
- **Jam signal (collision enforcement)**: To make sure that all stations involved in the collision will detect collision. A pattern of 32 bits.
- Collision backoff and retransmission method (**Truncated Binary Exponential Backoff Algorithm, BEBA**):
  - $n$  : number of collisions experienced ( $n \leq 16$ )
  - $k$  :  $\text{Min}(n, 10)$  -- Truncation
  - $r$  : Random delay time (unit: slot time),  $0 \leq r < 2^k$

# CSMA/CD Collision Handling

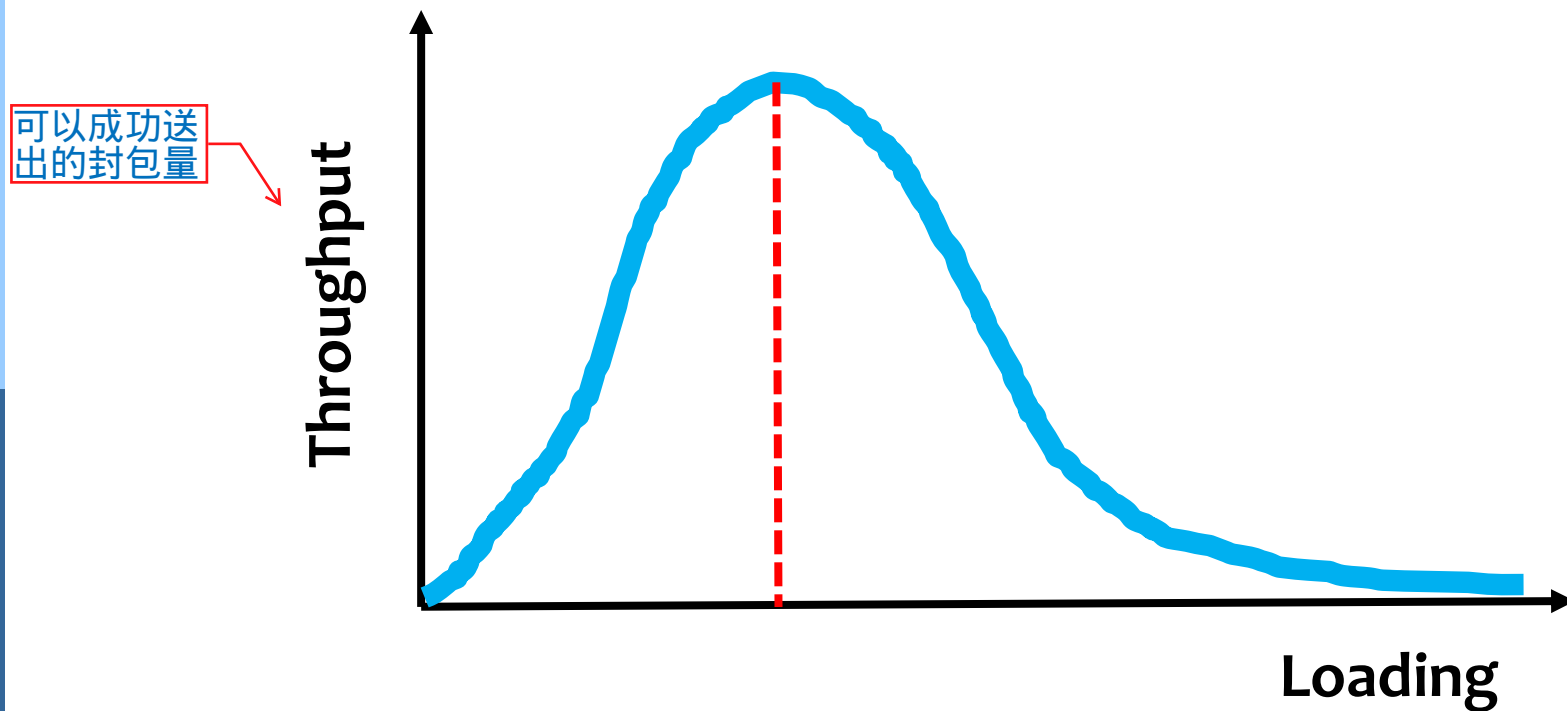
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- Slot time = 51.2 us.
- Disadvantage of BEBA:
  - **Last-in-First-out effect:** Stations with no or few collisions will have a better chance to transmit before stations that have waited longer.



# Ethernet Performance

- Ethernets work best under **lightly loaded** conditions.
- Under **heavy loads**, too much of the network's capacity is wasted by collisions.



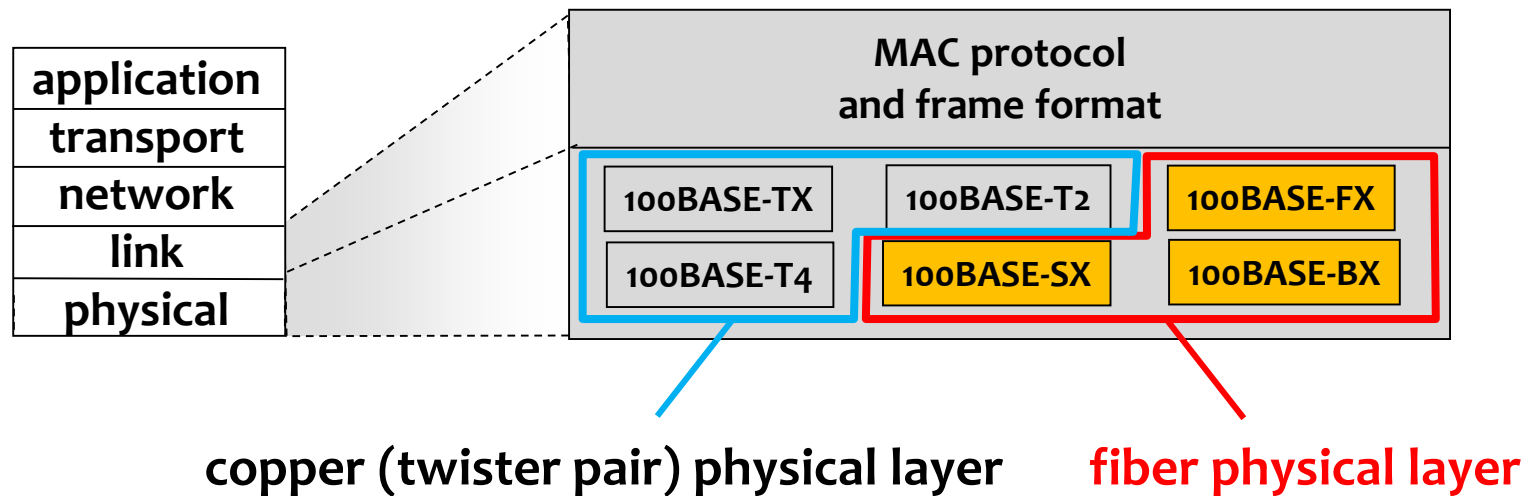
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# 802.3 Ethernet Standards: Link & Physical Layers

- **Many** different Ethernet standards
  - common MAC protocol (CSMA/CD) and frame format
  - different speeds: 2 Mbps, 10 Mbps, 100 Mbps, 1Gbps, 10Gbps, 100Gbps
  - different physical layer media: fiber, cable



# Summary

- MAC Protocol -- CSMA/CD
- Connection less, unreliable transmission
- Topology from **Bus** to **Star (switches)**
- Half-duplex transmission in Bus topology
  - Work best under **lightly loaded** conditions
  - Too much collision under **heavy load**
- Full-duplex transmission in Switch topology (point-to-point)
  - No more collisions !!
  - Excellent performance (wired speed)

努力送出，冲撞16次以内都会重送。但还有一种不可靠情况是，没有发生冲撞，但受到干扰，sender不知道，只要没有冲撞就不会重送，但receiver接收到后通过CRC检查发现封包错误，就会丢弃。所以不可靠

线路专用，还是跑CSMA/CD，但不会侦测到冲撞