

Introduction to Computer Networks



Foundation

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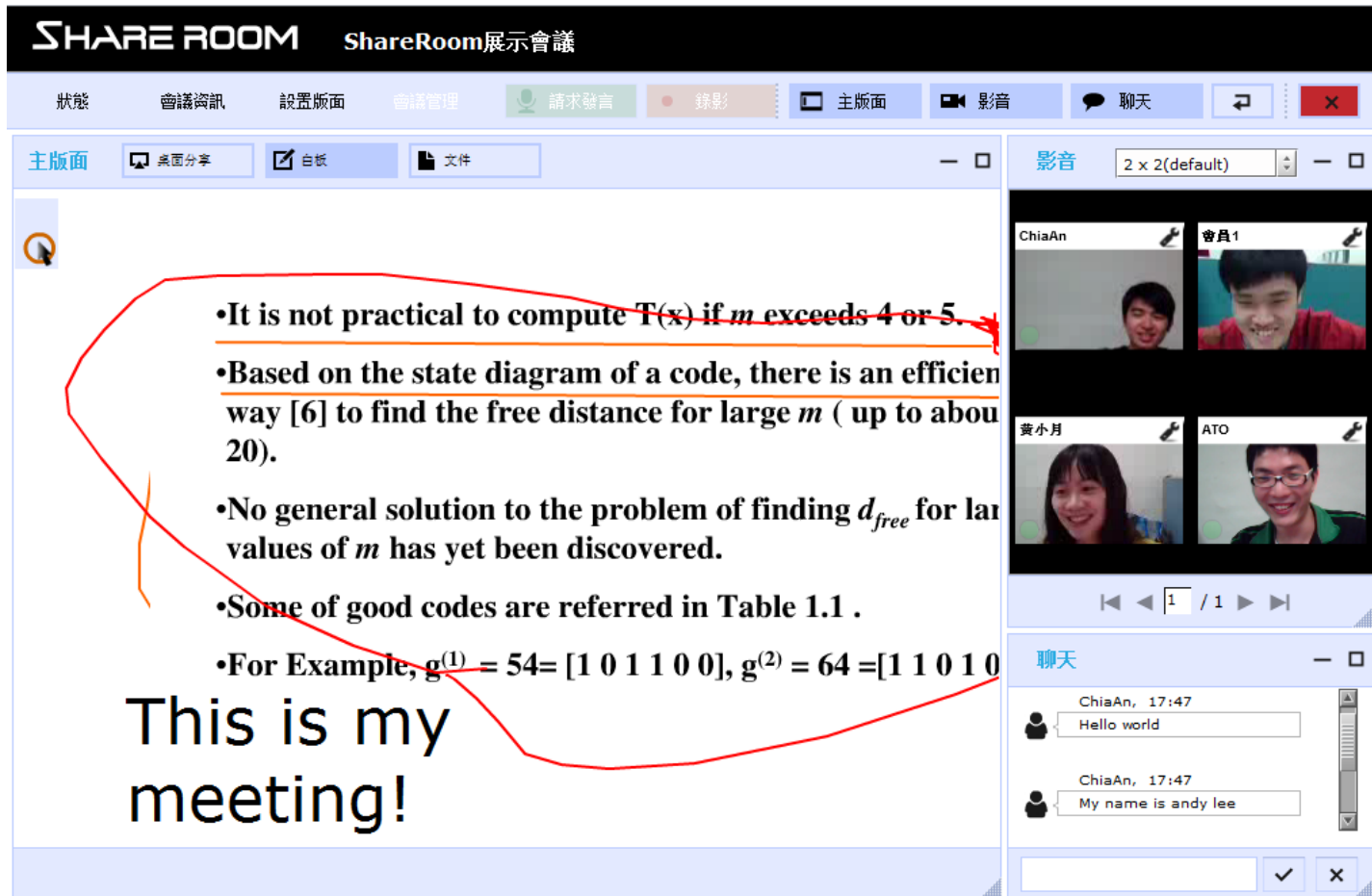
Outline

- **Applications**
- **Network Connectivity**
- **Network Architecture**
- **Network Performance**

Applications

- **Most people know about the Internet (a computer network) through applications**
 - **World Wide Web**
 - **On line games**
 - **Email (Gmail, hotmail,...)**
 - **Online Social Network (Facebook, twitter,...)**
 - **Streaming Audio Video (Youtube, ppstream, kkbox, ...)**
 - **File Sharing (dropbox, ...)**
 - **Instant Messaging (Skype, IM+, MSN, Line, WeChat,...)**
 - **...**

Example of an application



A multimedia application including video-conferencing

Application Protocol

■ URL

- Uniform resource locator
- <http://www.sharecourse.net/sharecourse/>

■ HTTP

- Hyper Text Transfer Protocol

■ TCP

- Transmission Control Protocol

■ 17 messages for one URL request

- 6 to find the IP (Internet Protocol) address
- 3 for connection establishment of TCP
- 4 for HTTP request and acknowledgement
 - ▶ Request: I got your request and I will send the data
 - ▶ Reply: Here is the data you requested; I got the data
- 4 messages for tearing down TCP connection

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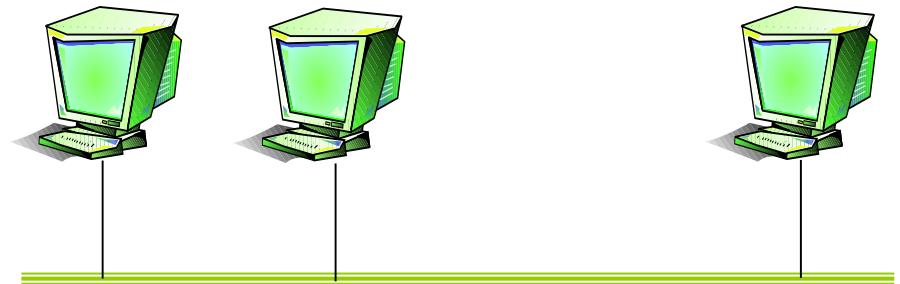
Network Connectivity

■ Important terminologies

- Link
- Nodes
- Point-to-point
- Multiple access
- Switched Network
 - ▶ Circuit Switched
 - ▶ Packet Switched
- Packet, message
- Store-and-forward



Point-to-point



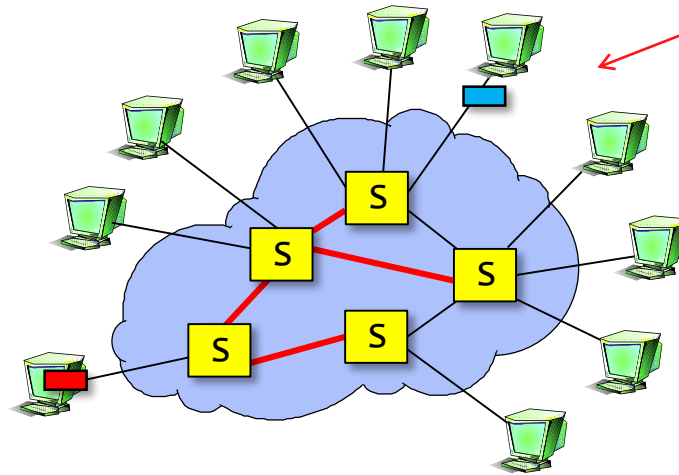
Multiple access

Network Connectivity

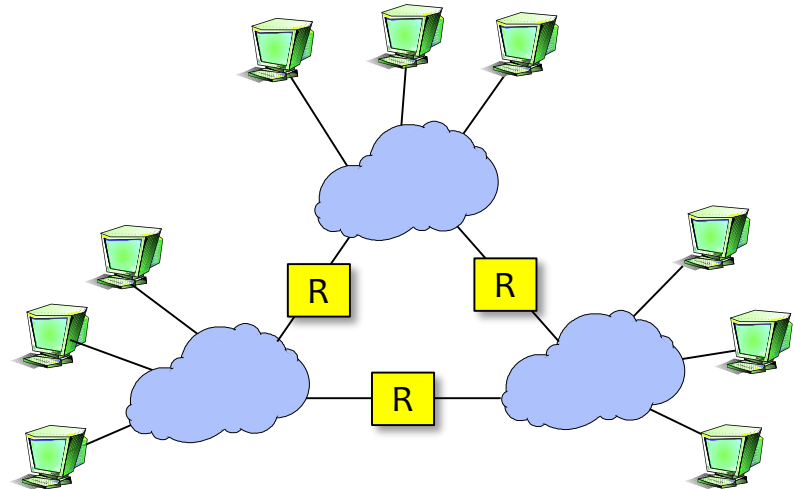
回路可以增加可靠性，避免连接中断造成影响；但可能造成重复收包等问题

■ Terminologies (contd.)

- Hosts
- Switches
- Spanning tree
- internetwork
- Router/gateway
- Host-to-host connectivity
- Address
- Routing
- Unicast/broadcast/multicast
- LAN (Local Area Networks)
- MAN (Metropolitan Area Networks)
- WAN (Wide Area Networks)

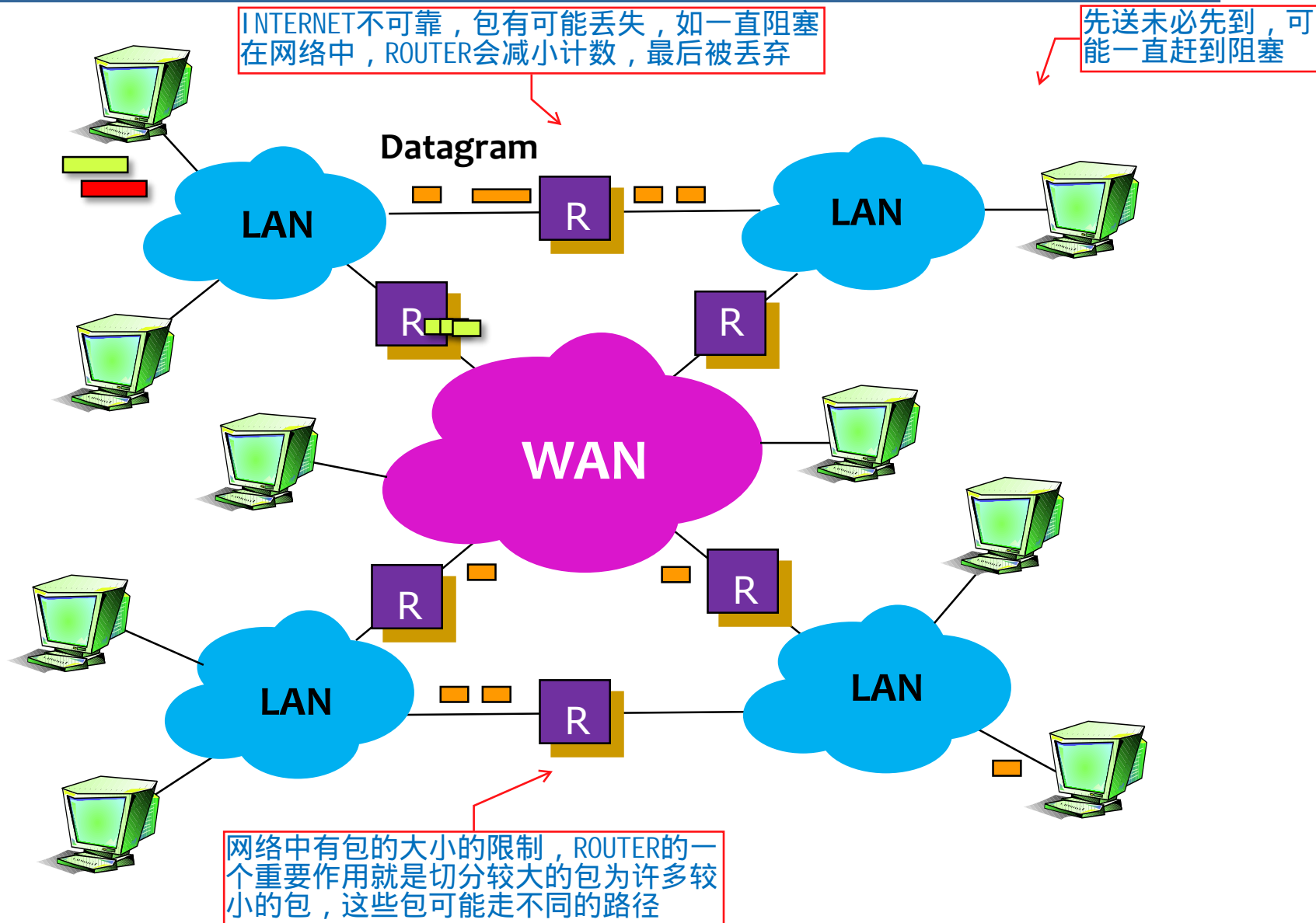


(a) A switched network



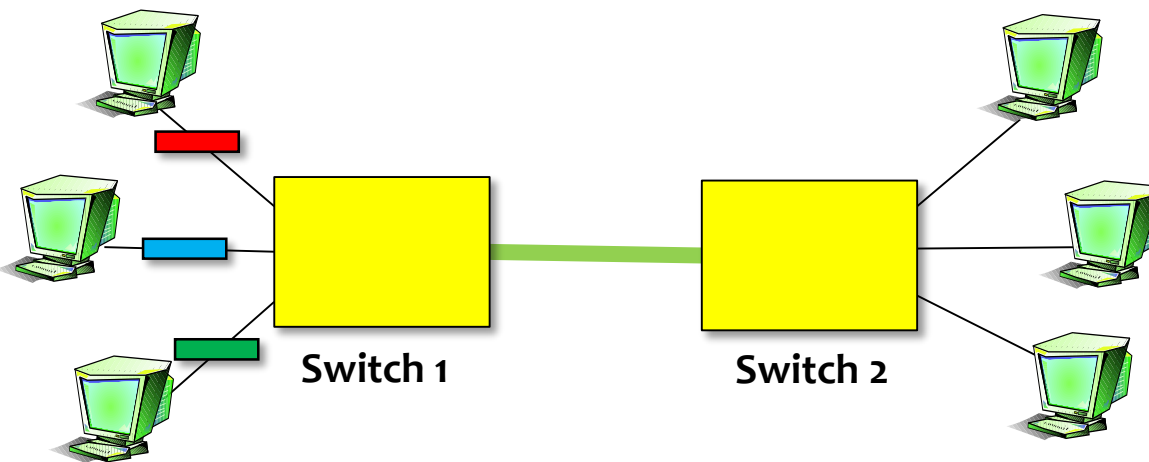
(b) Interconnection of networks

How datagrams are delivered in an Internet ?



Cost-Effective Resource Sharing

- Resource: links and nodes
- How to share a link ?
 - Multiplexing
 - De-multiplexing

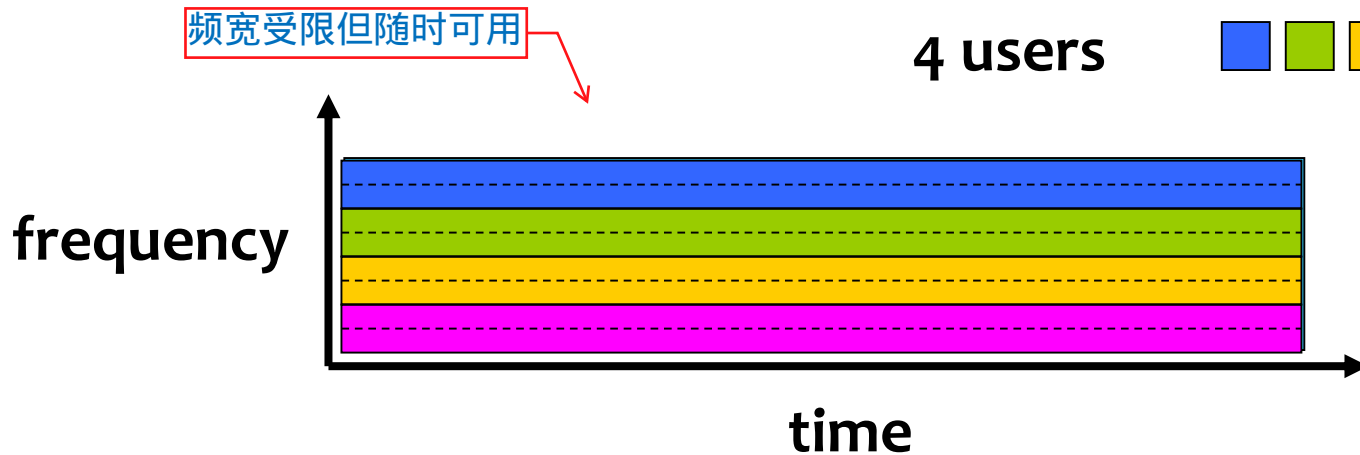


Multiplexing multiple logical flows over a single physical link

Cost-Effective Resource Sharing

■ FDM: Frequency Division Multiplexing

FDM



Example:

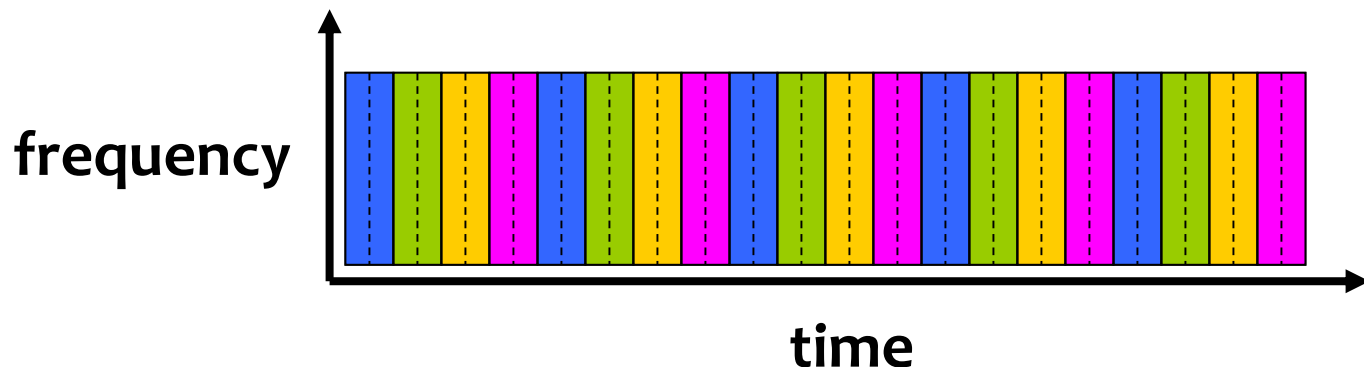
4 users



■ Synchronous Time-division Multiplexing (TDM)

- ▶ Time slots/data transmitted in predetermined slots

TDM



Cost-Effective Resource Sharing

■ Statistical Multiplexing

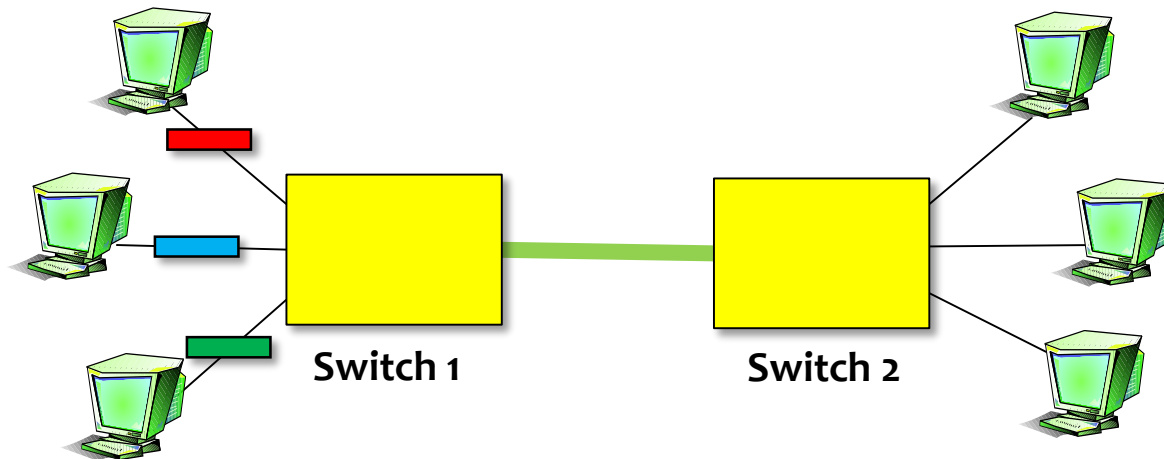
- Data is transmitted based on demand of each flow.
- What is a **flow**?
- Packets vs. Messages
- FIFO, Round-Robin, Priorities (Quality-of-Service (QoS))
- Congested ?

流量，弹性定义

共同占用一个队列，负载均衡时平均使用

每个主机各有一个队列

阻塞会导致丢包



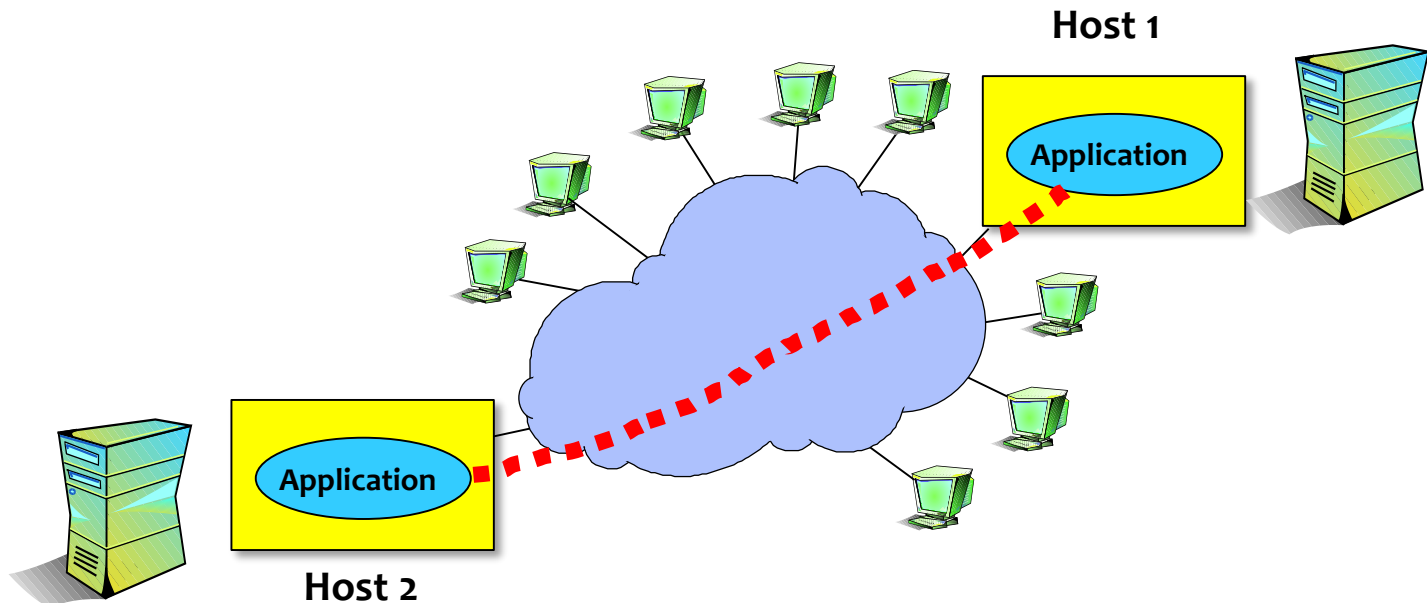
A switch multiplexing packets from multiple sources onto one shared link

Logical Channels

并非真正的线路，只是关系

■ Logical Channels

- Application-to-Application communication path or a pipe



Process communicating over an abstract channel

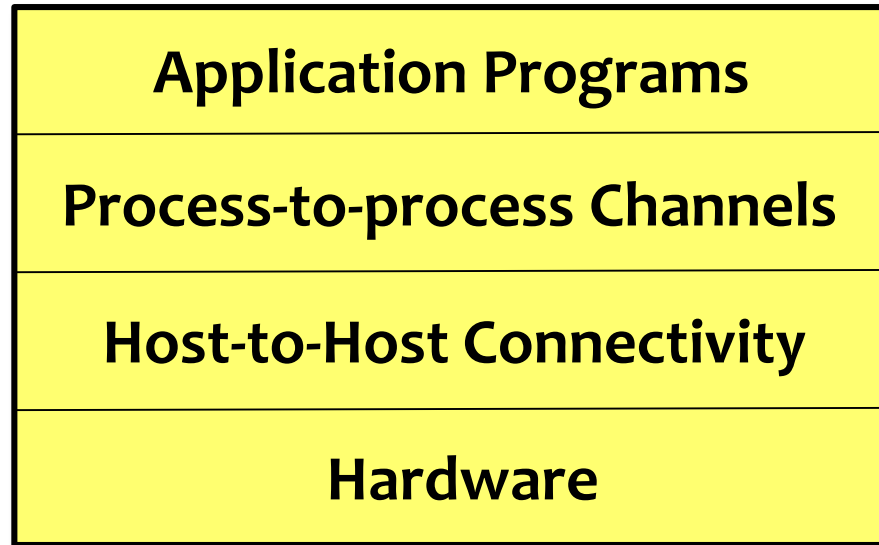
Network Reliability

- Network should **hide the errors**
- Bits are lost
 - Bit errors (1 to a 0, and vice versa)
 - Burst errors – several consecutive errors
- Packets are lost (Congestion)
- Links and Node failures
- Messages are delayed
- Messages are delivered out-of-order
- Third parties eavesdrop

Outline

- Applications
- Network Connectivity
- **Network Architecture**
- Network Performance

Network Architecture

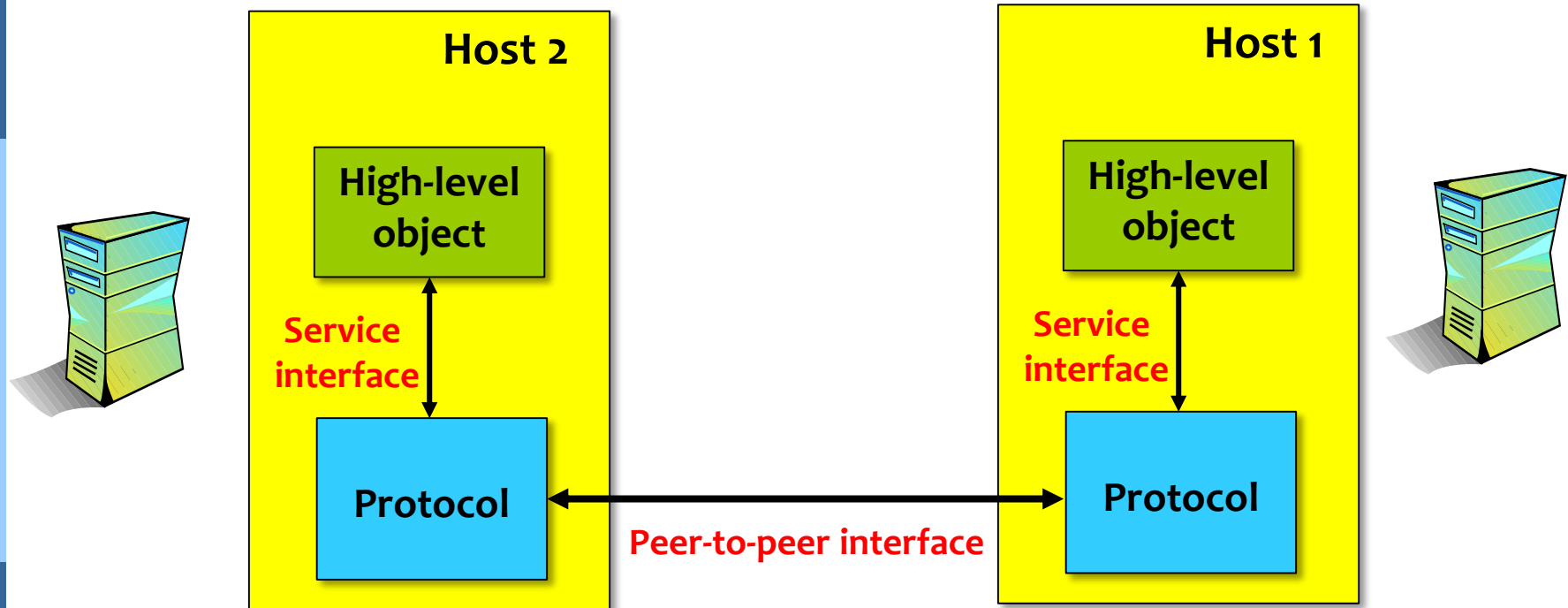


Example of a layered network system

Protocols

- **Protocol** defines the **interfaces** between
 - the layers in the same system and with
 - the layers of peer system
- Building blocks of a **network architecture**
- Each protocol object has two different interfaces
 - **Service interface**: operations on this protocol
 - **Peer-to-peer interface**: messages exchanged with peer

Protocol Interfaces

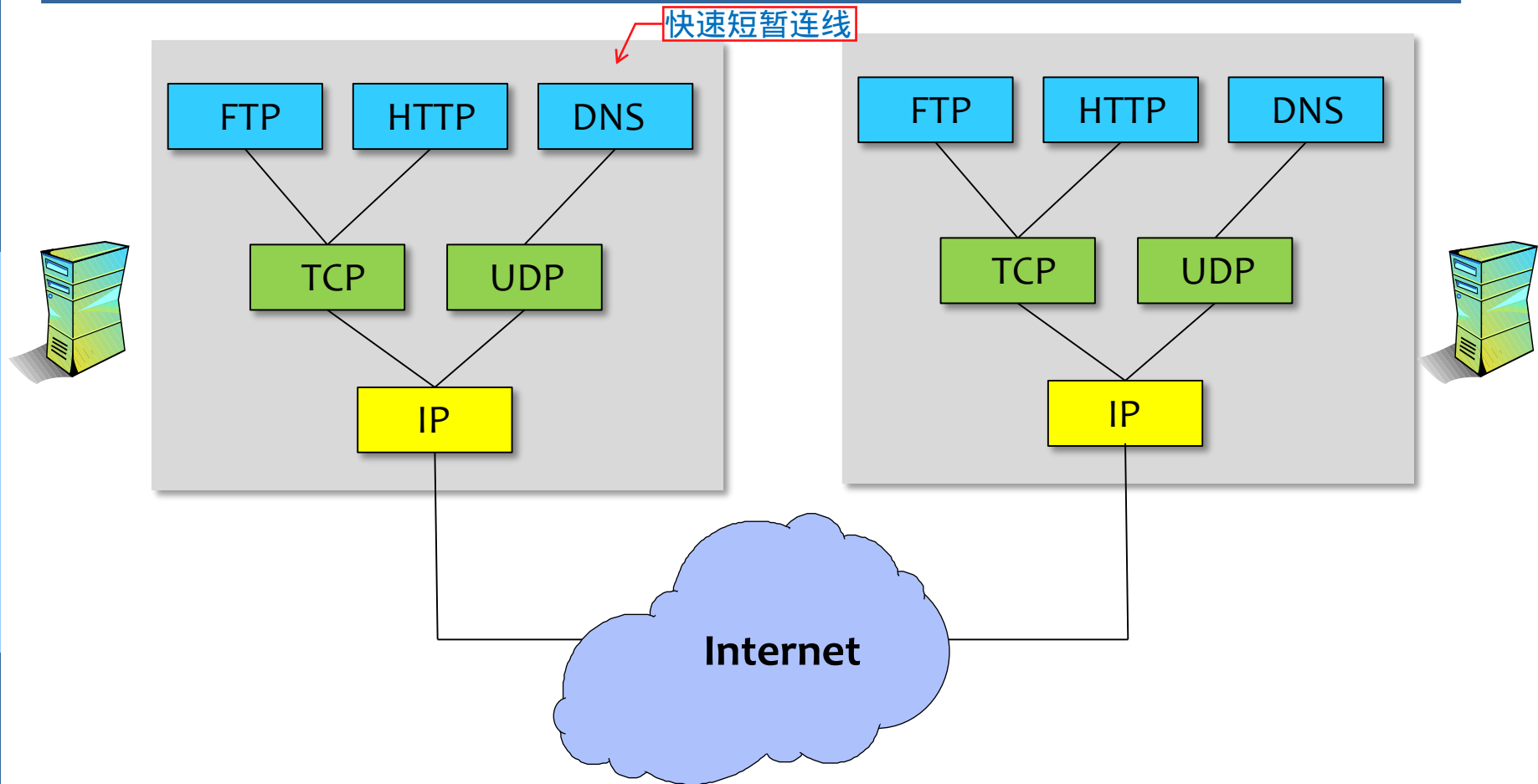


Service and Peer Interfaces for a protocol

Protocols

- **Protocol Specification:** pseudo-code, state transition diagram, message format
- **Interoperable:** when two or more protocols that implement the specification accurately
- **IETF: Internet Engineering Task Force**
 - Define Internet standard protocols

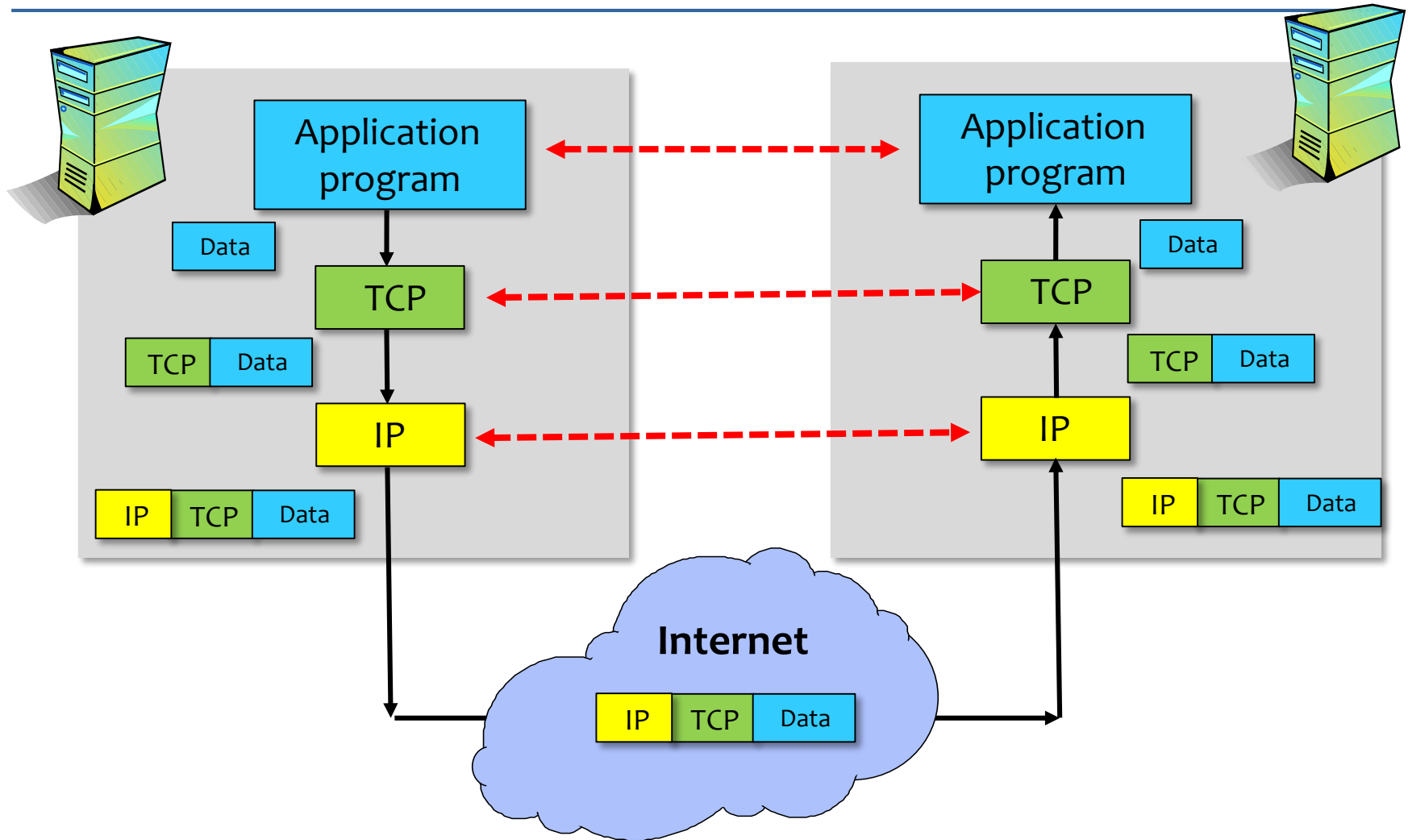
Protocol Architecture



Example of a protocol architecture

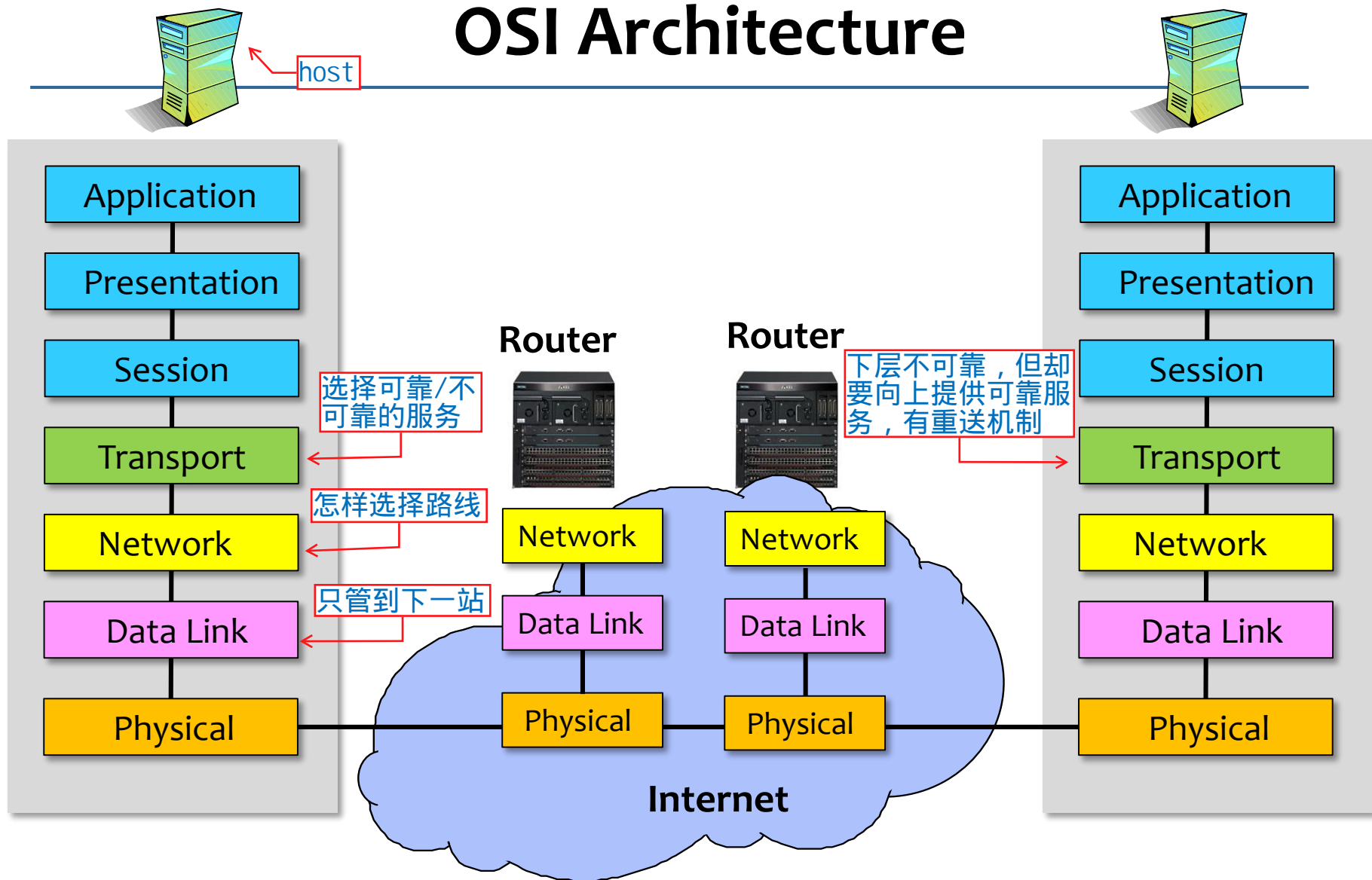
nodes are the protocols and links the “depends-on” relation

Encapsulation



High-level messages are encapsulated inside of low-level messages

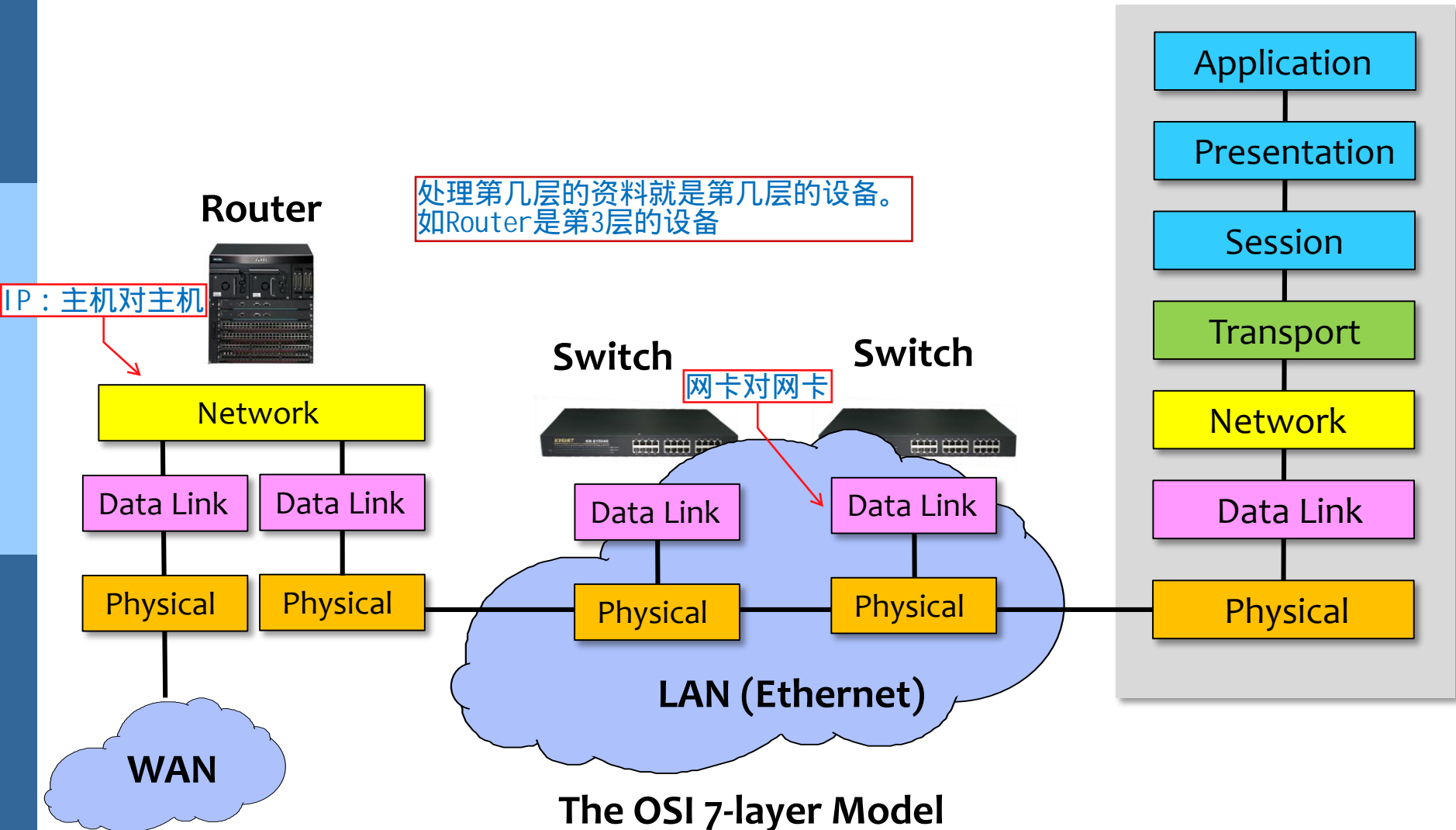
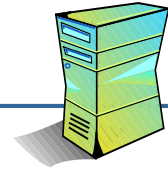
OSI Architecture



The OSI 7-layer Model

OSI – Open Systems Interconnection

OSI Architecture

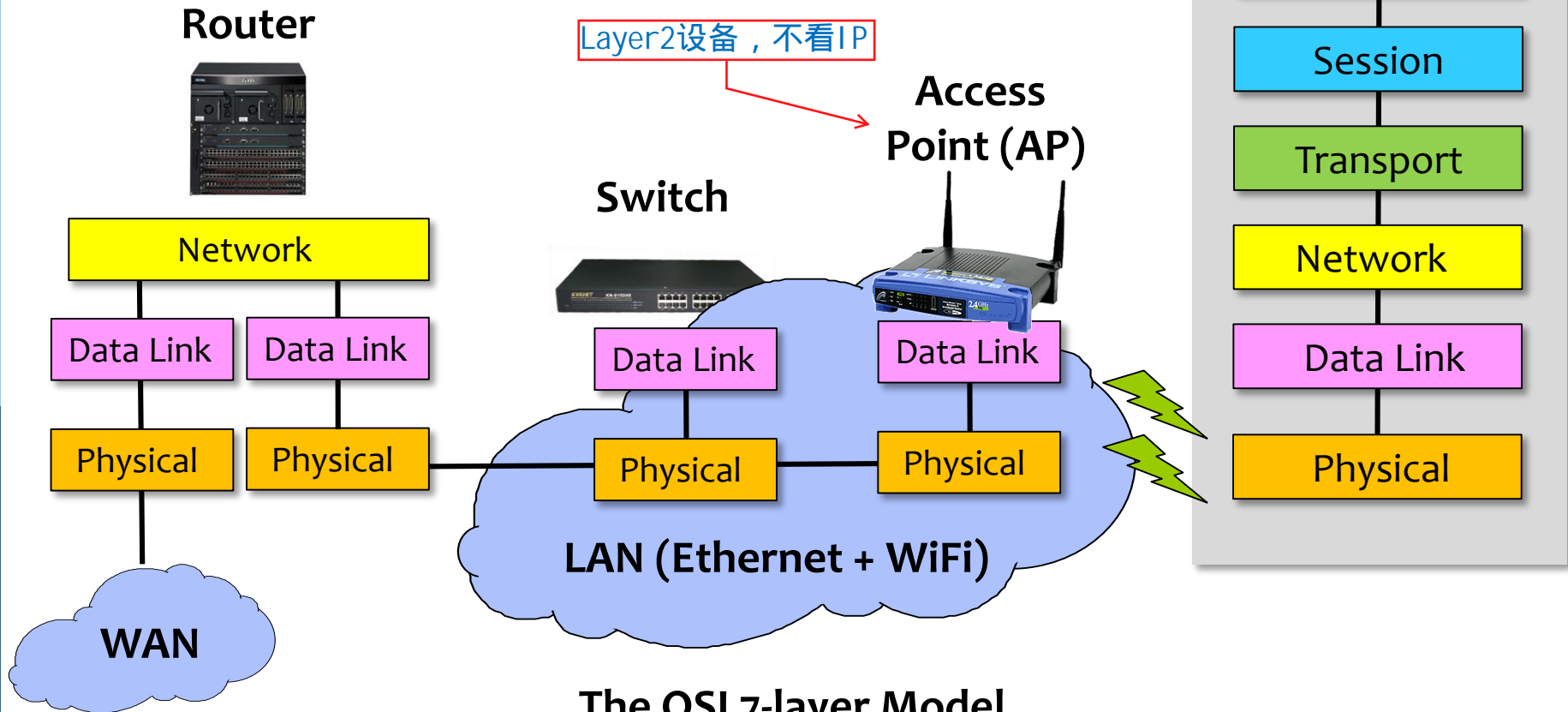
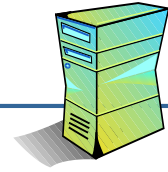


处理第几层的资料就是第几层的设备。
如Router是第3层的设备

网卡对网卡

The OSI 7-layer Model
OSI – Open Systems Interconnection

OSI Architecture



The OSI 7-layer Model

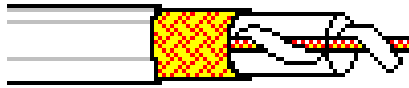
OSI – Open Systems Interconnection

Description of Layers

■ Physical Layer (如何將原始資料在 link 上傳輸)

- Handles the transmission of **raw bits over a communication link**

- ▶ Coaxial cable



- ▶ Twisted pair



- ▶ Optical Fiber

- ▶ Air space (wireless radio channel)



- Different Signal **Coding** schemes

发送方与接收方需
采用同样的clock



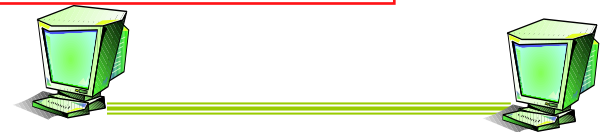
Description of Layers

■ Data Link Layer (如何將 frame 傳給直接相連的主機或設備)

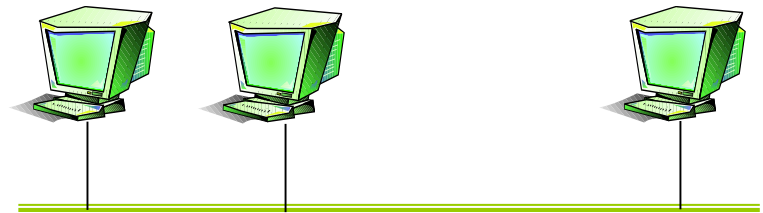
- Collects a stream of bits into a **frame**
- How to transmit a frame to a directly connected host (destination) ?
- MAC (Media Access Control Protocol) 凡是第二层的协议都是MAC

collision detection

- ▶ CSMA/CD (IEEE 802.3 Ethernet)
- ▶ CSMA/CA (IEEE 802.11 Wireless LAN)

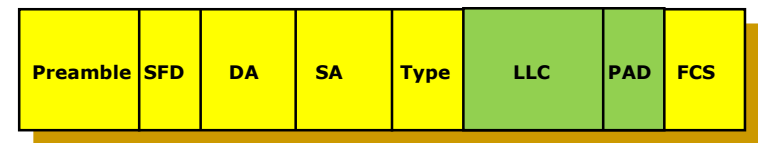
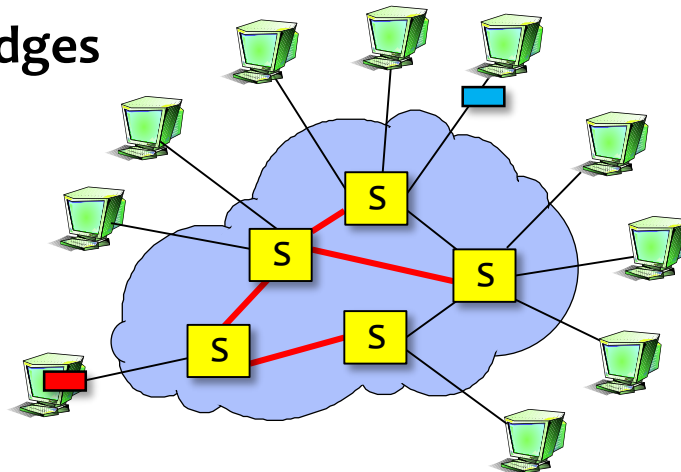


Point-to-point



Multiple access

- Layer 2 devices
- Switches
- Bridges



加入错误检查

Description of Layers

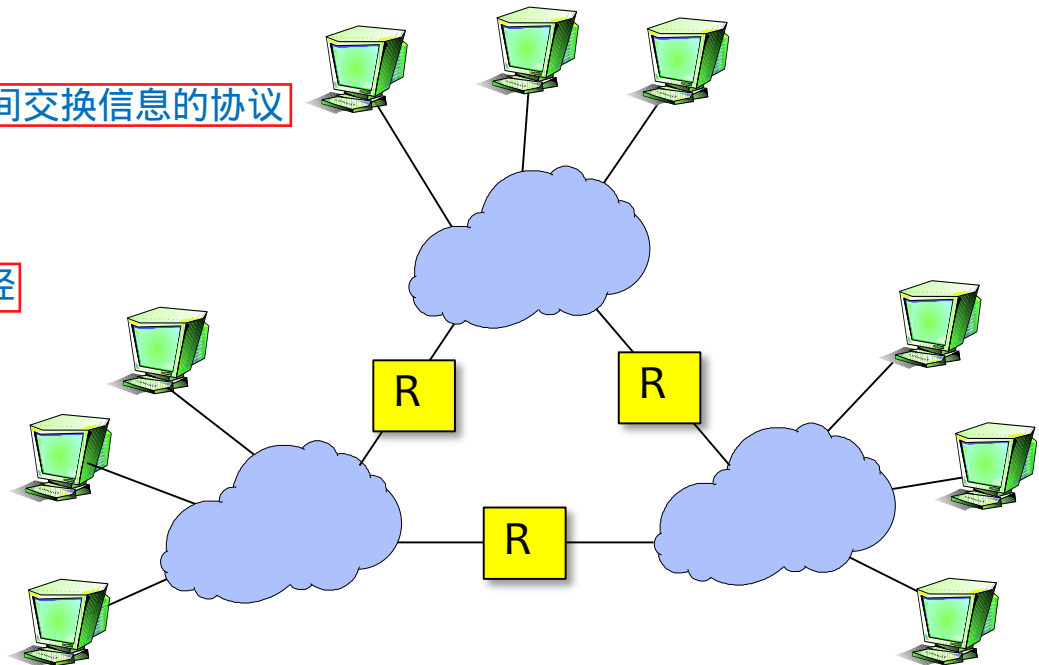
■ Network Layer (如何將封包透過 Internet 送給目的地主機)

- How to transmit frames to a host via the Internet ?
- Handles **routing** among nodes within a packet-switched network
- Data exchanged between nodes in this layer is called a **packet**
- *IP protocol*
- **Routers**
- **Routing protocols**
 - ▶ **RIP**
 - ▶ **OSPF**
 - ▶ **BGP**
- **Routing Tables**

查表：动态变更

Router之间交换信息的协议

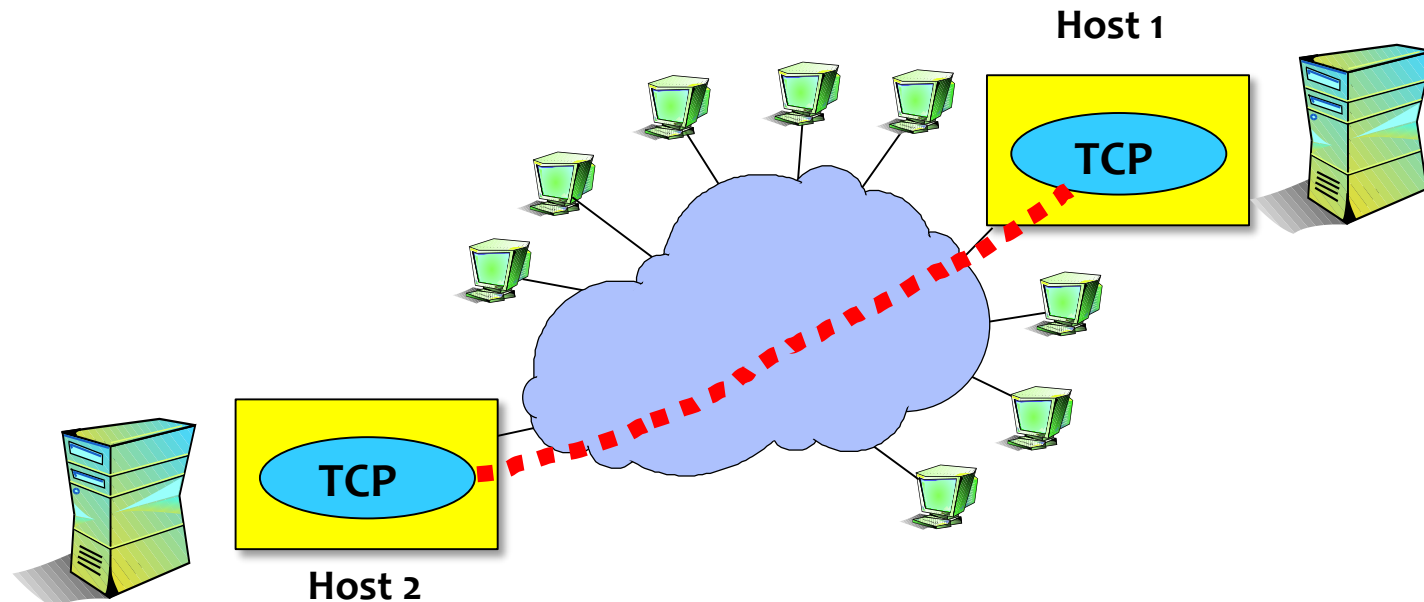
计算最短路径



Interconnection of networks

Description of Layers

- Transport Layer (提供不同主機 processes 之間的資料傳送)
 - Implements a **process-to-process** channel
 - Unit of data exchanges in this layer is called a **message**
 - **TCP (Transmission Control Protocol)** – Reliable service
 - **UDP (User Datagram Protocol)** – Unreliable service



Description of Layers

■ Session Layer

- Provides a name space that is used **to tie together the potentially different transport streams** that are part of a single application

■ Presentation Layer

- Concerned about the **format of data** exchanged between peers

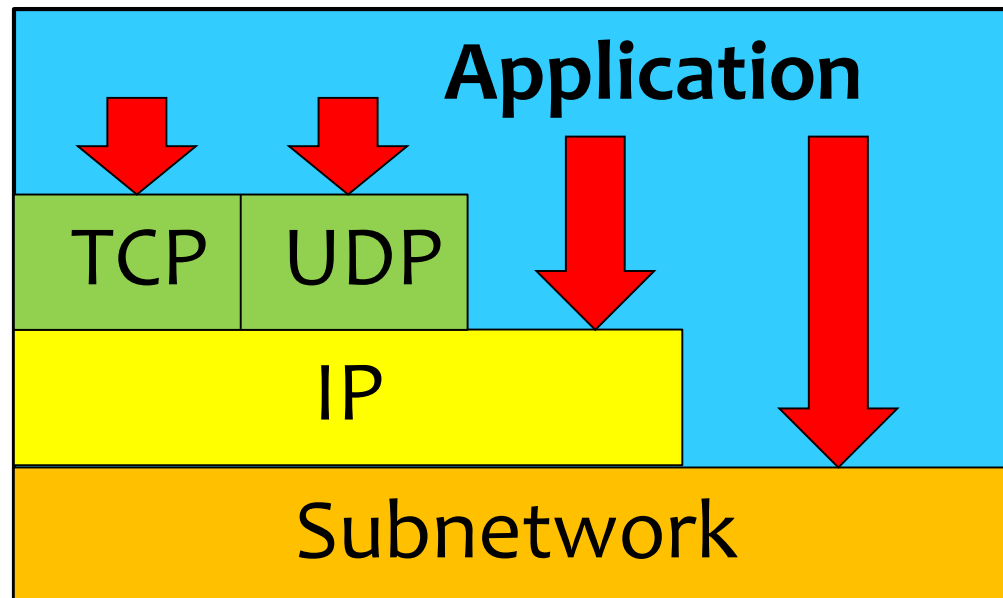
■ Application Layer

- Standardize common type of exchanges
- FTP/E-mail/DNS/HTTP/Browsers/FB,

The transport layer and the higher layers typically run only on end-hosts and not on the intermediate switches and routers

Internet Architecture

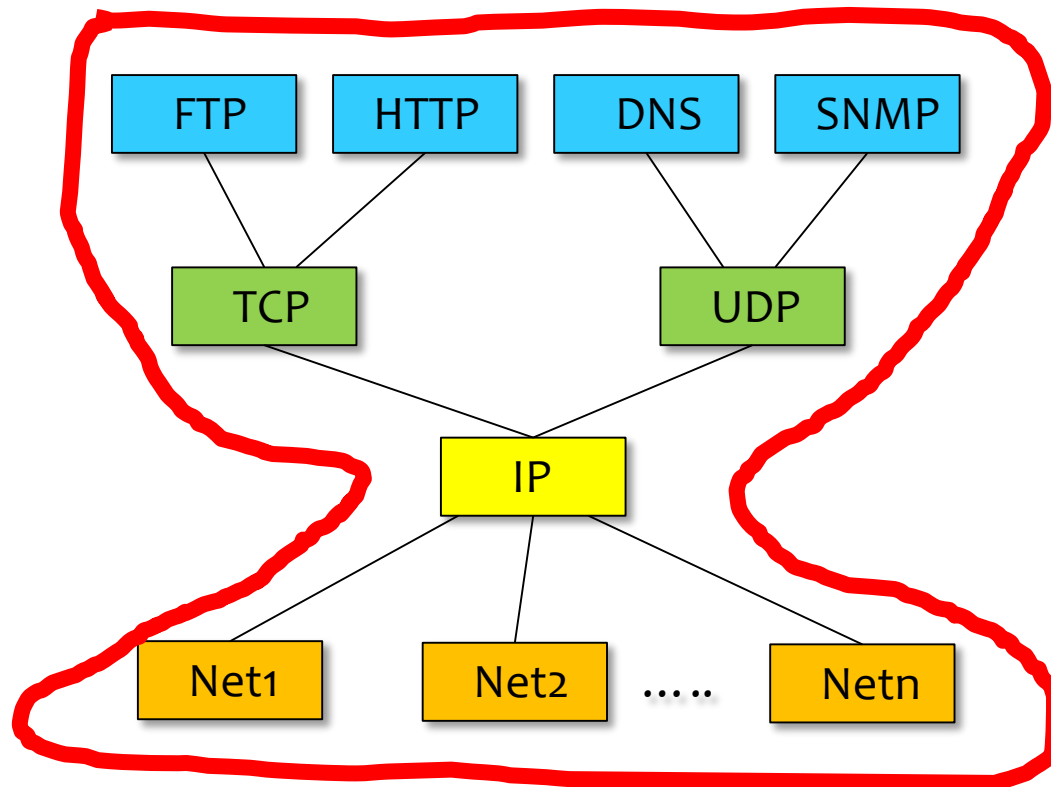
- Defined by IETF
- Three main features
 - **Does not imply strict layering.** The application is **free to bypass** the defined transport layers and to directly use IP or other underlying networks



Internet Architecture

沙漏状

- **An hour-glass shape** – wide at the top, narrow in the middle and wide at the bottom. IP serves as the focal point for the architecture



Internet Architecture

- In order for a new protocol to be officially included in the architecture, there needs to be both **a protocol specification** and **at least one** (and preferably two) **representative implementations** of the specification

Outline

- Applications
- Network Connectivity
- Network Architecture
- **Network Performance**

Network Performance

■ Bandwidth

- Width of the **frequency band**
- **Number of bits per second** that can be transmitted over a communication link

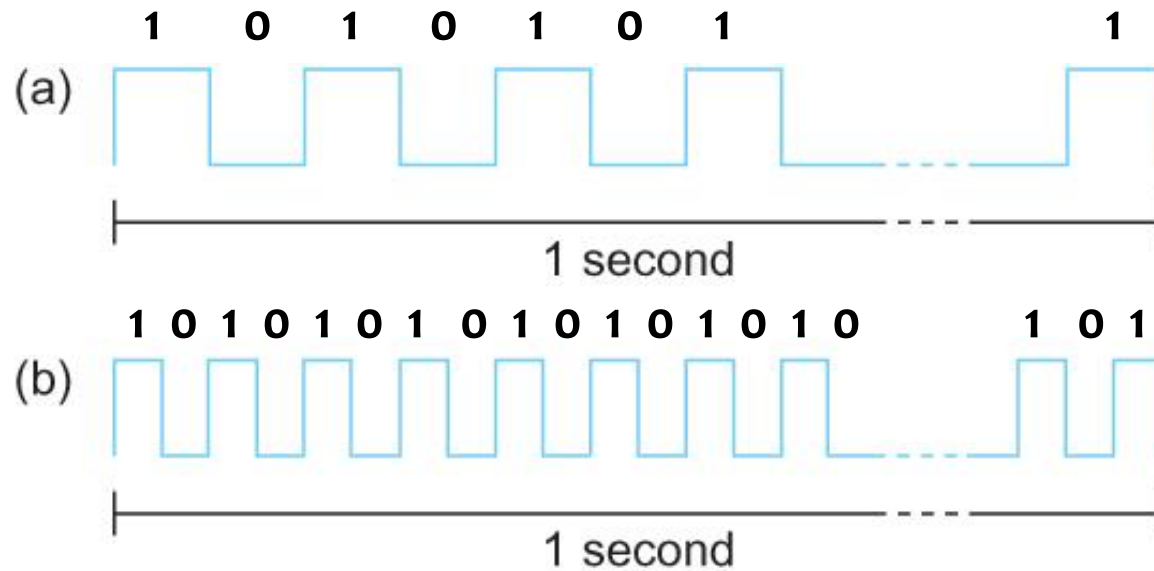
■ 1 Mbps: 1×10^6 bits/second

■ 1×10^{-6} seconds to transmit each bit or imagine that a timeline, now each bit occupies 1 micro second space.

■ On a 2 Mbps link the width is 0.5 micro second.

■ Smaller the width more will be transmission per unit time.

Bandwidth



Bits transmitted at a particular bandwidth can be regarded as having some width:

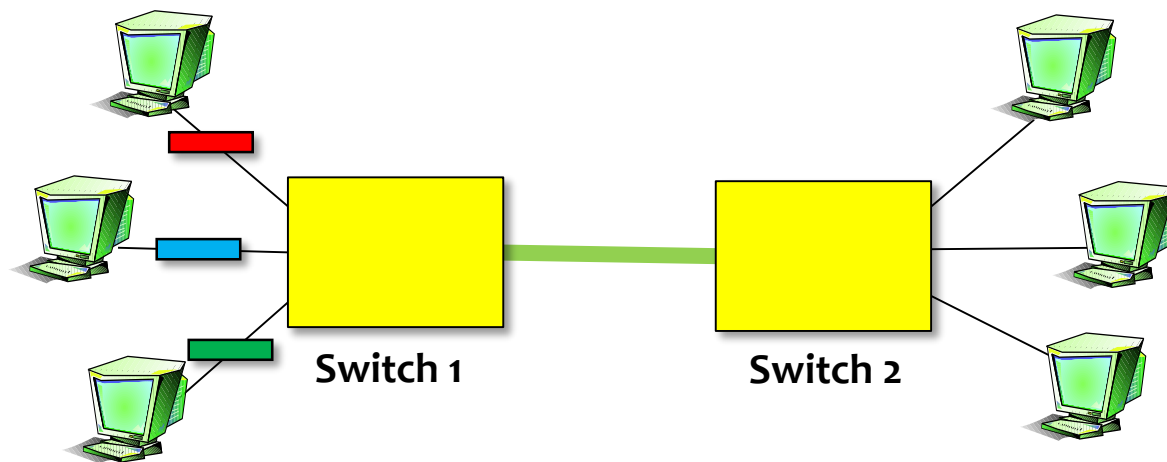
(a) bits transmitted at 1Mbps (each bit 1 μ s wide);

(b) bits transmitted at 2Mbps (each bit 0.5 μ s wide).

越高速越难正确接收

Network Performance

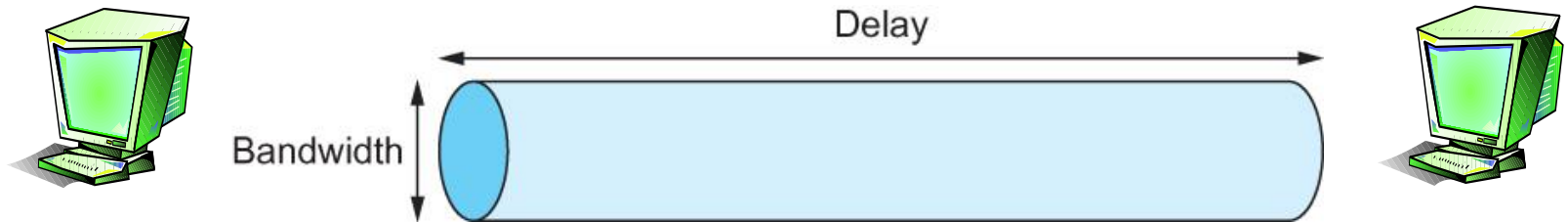
- **Latency** = Propagation time + transmission time + queuing time
- **Propagation** time = distance/speed of light
- **Transmission time** = size/bandwidth



- One bit transmission => propagation is important (短資料很快就送完, 但需要長時間才能傳到對方, 資料已送完, 但前導資料還未到達對方)
 - Propagation time >> transmission time
- Large bytes transmission => bandwidth is important (長資料很慢才能送完, 未送完前, 前導資料已到對方)
 - Transmission time >> propagation time

Delay X Bandwidth

- The channel between a pair of processes can be viewed as a pipe
- **Latency (delay): length** of the pipe
- **Bandwidth: width** of the pipe
- Delay x Bandwidth means how many data can be stored in the pipe
- For example, delay of 80 ms and bandwidth of 100 Mbps
 - ⇒ 80×10^{-3} seconds $\times 100 \times 10^6$ bits/second
 - ⇒ 8×10^6 bits = 8 M bits = 1 MB data.



Network as a pipe

Delay X Bandwidth

- Relative importance of **bandwidth** and **latency** depends on application
 - For large file transfer, bandwidth is critical
 - For small messages (HTTP, NFS, etc.), latency is critical
 - Variance in latency (**jitter**) can also affect some applications (e.g., audio/video conferencing)

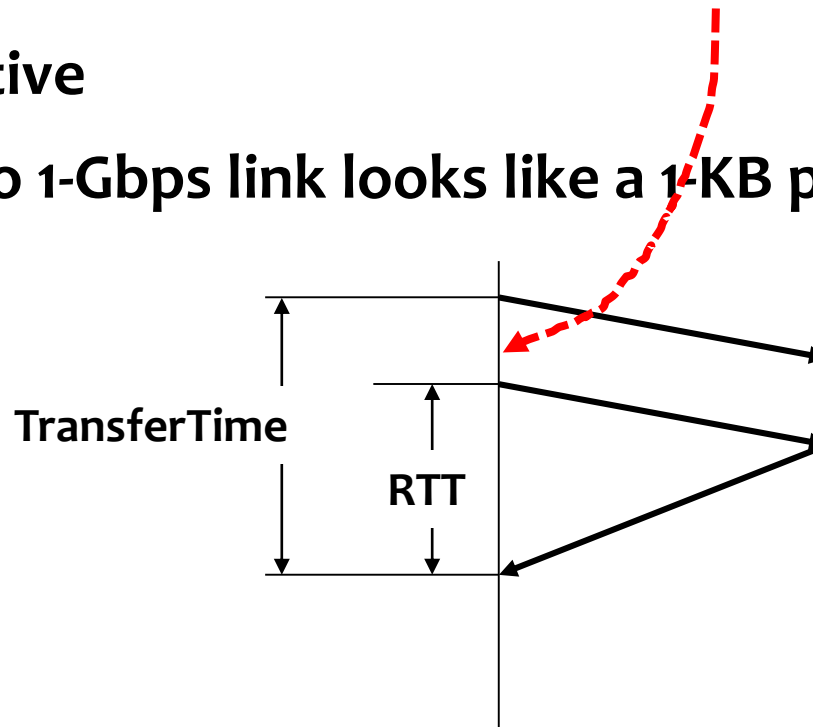
Delay X Bandwidth

- if the sender keeps the pipe full, **delay x bandwidth** is the number of bits the sender must transmit before the first bit arrives at the receiver
- Takes another one-way latency to receive a response from the receiver
- The sender will not fully utilize the network if the sender does not fill the pipe
 - ▶ send a whole delay x bandwidth product's worth of data before it stops to wait for a signal
 - ▶ 在停下來等對方回應之前應該要傳送 delay x bandwidth 的資料量

Throughput

■ Infinite bandwidth

- RTT (Round Trip Time) dominates
- **Throughput** = $\text{TransferSize} / \text{TransferTime}$
- $\text{TransferTime} = \text{RTT} + \text{TransferSize} / \text{Bandwidth}$
- Its all relative
- 1-MB file to 1-Gbps link looks like a 1-KB packet to 1-Mbps link



Summary

- A **layered architecture** for computer network
 - Physical Layer
 - Data Link Layer
 - Network Layer
 - Transport Layer
 - Session layer / Presentation Layer /Application layer
- **Two performance metrics** used to analyze the performance of computer networks
 - Bandwidth
 - Delay