Encapsulation

Programming Languages
CS 214



Review

A type consists of two things:

- •a set of ____; and
- •a set of onto those values.

An abstract data type (ADT) consists of:

- a collection of ____: $type_1 \times type_2 \times ... \times type_n$
- •a collection of on the data:

$$F(type_1 \times type_2 \times ... \times type_n) \rightarrow ADT$$

 $G(ADT) \rightarrow type_1$

$$H(ADT) \rightarrow \emptyset$$

Obviously, these two are related...



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Encapsulation

An encapsul	lation mechanism is a language's construct for into
a single syr	tactic structure.
Two differen	nt encapsulation mechanisms have evolved:
	, a mechanism that lets programmers create s that encapsulate data and operations; and
	(aka <i>package</i>), a mechanism that lets mers store new types and their operations in a ontainer.

The evolutionary history of these provides useful context, so we'll examine the history of each separately...



ADTs

In the early 1970s, *imperative* programming languages had evolved to the point where much of programming was building _______, which consisted of an ______ and its supported ______.

For example, a Stack ADT consists of:

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- *Initialization*(&Stack) $\rightarrow \emptyset$
- $isEmpty(Stack) \rightarrow bool$
- $isFull(Stack) \rightarrow bool$

- $push(value \times \&Stack) \rightarrow \emptyset$
- $pop(\&Stack) \rightarrow \emptyset$
- $top(Stack) \rightarrow value$

The ADT's operations make up its _____ through which users are to interact with the ADT.

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ADT Example in C

```
/* IntStack.h (minus precondition checks)*/
#define STACK MAX 32
                                        top
typedef struct StackStruct {
  int myTop;
                                      values
  int myValues[STACK_MAX];
                                            [0] [1] [2] [3] ... [31]
} IntStack;
void init(IntStack* sRef) { sRef->myTop = -1; }
int isEmpty(IntStack s) { return s.myTop < 0; }</pre>
int isFull(IntStack s) {return s.myTop >= STACK MAX-1;}
void push(int value, IntStack* sRef) {
  sRef->myTop++;
  sRef->myValues[sRef->myTop] = value;
void pop(IntStack* sRef) { sRef->myTop--; }
int top(IntStack s) { return s.myValues[s.myTop]; }
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```



Problem 1

Nothing prevents a programmer from

Instead of writing:

a programmer can write:

If we upgrade our array-based stack to a linked stack:

```
while (!isFull(s))
 // ... do something with s
while ( s.myTop < STACK_MAX )</pre>
 // ... do something with s
/* IntStack.h ... */
typedef struct Node {
  int value;
  struct Node * next;
} IntStackNode;
typedef struct StackStruct {
   IntStackNode * top;
} IntStack;
```

then the programmer's code (\(^\text{maintenance costs}\)!



Problem 2

Nothing prevents the programmer from #include "IntStack.h" int peekUnderTop(IntStack s) return s.values[s.top-1]; Such operations will also ______ if we change the implementation details that underlie the *Stack* ADT (↑ maintenance costs)! Recall: High maintenance costs led industry from *spaghetti coding* to structured programming. Eliminating the maintenance costs of "broken" code led industry from structured programming to _____ programming.



Modules and Packages

The problem with structured programming was that it did	
nothing to	•

In 1977, Wirth designed *Modula* with a new "container" construct in which a type and its operations could be stored.

- Wirth called this container the _____.
- Rather than thinking of a type as values and operations, Wirth considered a type to be *just values* (i.e., data).
- The *module* was Wirth's construct for "wrapping" a type and its operations together (i.e., building an ADT).
- Fortran (90 and later) also provides a *Module* construct

 In the 1980s, Ada adopted a similar approach for ADTs, but called their container the ______ instead of the module.

ADT Example in Ada

```
-- IntStackPackage.ads is the IntStackPackage specification
  the Ada equivalent of a C header file
package IntStackPackage is
  type IntStack is private;
  procedure init(s: in out IntStack);
  function is Empty(s: in IntStack) return Boolean;
  function isFull(s: in IntStack) return Boolean;
  procedure push(value: in Integer; s: in out IntStack);
  procedure pop(s: in out IntStack);
  function top(s: in IntStack) return Integer;
                                             All declarations before
 private
  STACK_MAX: constant Integer := 32;
                                                private are visible
                                              externally; those after
  type IntStack is
   record
                                              private are local to the
    myTop: Integer;
    myValues: array(1...STACK MAX) of Integer;
                                                    package.
  end Stack;
end IntStackPackage;
```

ADT Example in Ada (ii)

```
-- IntStackPackage.adb is the IntStackPackage body
     the Ada equivalent of a C implementation file
package body IntStackPackage is
  procedure init(s: in out IntStack) is
  begin
      s.myTop:= 0;
  end init;
  function isEmpty(s: in IntStack) return Boolean is
  begin
      return s.myTop < 1;</pre>
  end isEmpty;
  procedure push (value: in Integer; s: in out IntStack) is
  begin
      s.myTop:= s.myTop + 1; s.myValues(s.myTop):= value;
  end;
   -- ... definitions of isFull(), pop(), top(), ...
end IntStackPackage;
```

Package Specifications

The "public section" of the specification	creates its
 Nothing else in the package is accessible 	
 If a programmer wishes to use the ADT, the declarations in its interface. 	they must do so using
An Ada package specification differs fr	om a C header file:
 Its private section allows the package to _ 	
from the programmer (everything in	n a header file is public).
- The specification file <i>must be compiled</i> be (and before the package body can be comp	
By separating the ADT's public	from its private
, a programmer can	nnot write programs
that depend upon the ADT's implemente	ation details.

Example Usage

Given such an ADT, a programmer can write:

```
-- IntStackTest.adb
with TextIO, IntStackPackage; use TextIO, IntStackPackage;
procedure IntStackTest is
  s1, s2: IntStack;
  i: Integer;
begin
  init(s1);
  while (not isFull(s1)) loop
    get(i); push(i, s1);
  end loop;
  s2 := s1;
  while (not isEmpty(s2)) loop
    put (top(s2)); pop(s2);
  end loop;
end IntStackTest;
```

In modular programming, ADT operations are subprograms that

Used in this way, a package/ module is a container in which an ADT can be "wrapped".

Modules/Packages As "Objects"

Modules/packages can also be used as "objects":

```
-- IntStack.ads is the IntStack specification
package IntStack is
  function isEmpty() return Boolean;
  function isFull() return Boolean;
  procedure push(value: in Integer);
  procedure pop;
  function top() return Integer;
end IntStack;
```

All of these identifiers are public (there is no private section).

The operations do *not* receive the ADT via a parameter.

There is no *init* subprogram in this kind of module/package (we'll see why shortly).



"Object" Package Bodies

```
-- IntStack.adb is the IntStack body
package body IntStack is
 STACK_MAX: constant Integer := 32;
 myTop: Integer;
 myValues: array(1..STACK_MAX) of Integer;
 function isEmpty() return Boolean is
 begin
   return myTop < 1;
 end;
 procedure push (value: in Integer) is
 begin
   myTop:= myTop+1; myValues(myTop):= value;
 end push;
 -- ... definitions of isFull, pop, top, ...
begin
  myTop:= 0;
end IntStack;
```

Note 1: All of the implementation details are here in the body, making them

Note 2: A package may have an

at its end that is executed when a program using the package is run...

"Object" Modules in Use

Such modules/packages can be used in an object-like way:

```
-- IntStackTest.adb tests the IntStack package
with Text IO, IntStack; use Text IO;
                                         In this approach, a
                                          module/package
procedure IntStackTest is
  i: Integer;
                                          superficially resembles an
begin
                                          OO-language object
  while (not IntStack.isFull()) loop
                                          (created from a class)...
    get(i);
    IntStack.push(i);
                                         But a module/package is ____
  end loop;
                                                     _, so it cannot
  while (not IntStack.isEmpty()) loop
                                          be used to create variables.
    put(IntStack.top());
    IntStack.pop;
                                         → Only one such "object"
  end loop;
                                          can exist at a time.
end IntStackTest;
```

Generic Packages

```
Ada also allows packages to be given _
providing a way to circumvent the "one object" problem:
-- Stack.ads is the generic Stack specification
generic
  type Item is private; -- Item is a type parameter
                               -- size is a data parameter
  Integer size;
package Stack is
  function isEmpty() return boolean;
  function isFull() return boolean;
 procedure push(v: in Item);
 procedure pop();
 function top() return Item;
end Stack;
Ada's keyword _____ tells the compiler that the parameters Item
 and size will be supplied by the ADT's user (ideal for containers).
Such a Stack stores "generic" Items, instead of "hardwired" Integers.
```



Generic Package Bodies

```
-- Stack.adb is the generic Stack body
package body Stack is
 myCapacity: Integer := size;
 myTop: Integer;
 myValues: array(1..myCapacity) of Item;
 function isEmpty() return Boolean is
 begin
   return myTop < 1;</pre>
 end;
 procedure push (value: in Item) is
 begin
  myTop:= myTop+1; myValues(myTop):= value; Fortran's module
 end push;
 -- ... definitions of isFull, pop, top, ...
begin
  myTop := 0;
end Stack;
```

This far more elegant than ___ the compiler already knows that Stack is a generic package because the spec. is compiled first.

Recent versions of have added this generic mechanism.

Generic Instantiation

Now we can dynamically create multiple *Stack* "objects":

```
-- StackTest.adb tests the generic Stack package
with Text IO, Stack; use Text IO;
procedure StackTest is
  i: Integer;
  package intStack1 is new Stack(integer, 8);
  package intStack2 is new Stack(integer, 8);
begin
  while (not intStack1.isFull()) loop
                                             This permits generic
    get(i);
                                              Ada packages to be
    intStack1.push(i);
  end loop;
                                              constructed and
  intStack2:= intStack1;
                                              operated on in a way
                                              similar to objects.
  while (not intStack2.isEmpty()) loop
    put(intStack2.top());
    intStack2.pop;
  end loop;
end StackTest;
```



History: Simula

Back in 1967, Dahl & Nygaard not programming consists of	ed that at its simplest,
That is, variables are of values and operations, programm operating on variables.	, and since types consist ming can be reduced to
Dahl & Nygaard were working on or representation of "real world" objected real-world processes could be more	ects in software, so that
Their language was Simula (Simple	universal language) and

it provided useful Simulation and Process constructs...

Simula Classes

Dahl & Nygaard reasoned that if types are *values plus* operations, then a language should provide a syntactic structure that explicitly combines data and operations.

They took the _____ construct, extended it to store *subprograms* as well as *data* and christened it the _____ (from mathematics).

Subprograms (operations)

Data (state information)

Their *class* construct provided for creating ______, which could then be used to *declare variables*.

In time, variables were replaced by *objects* and objectoriented programming (OOP), culminating in

Example: An IntStack Class

```
! IntStack.sim (minus precondition checks);
Class IntStack(Size);
                                     ! Classes can have parameters;
  Integer Size;
                                     ! Params precede 'Begin';
Begin
                                     ! Attribute variables;
                                        Data encapsulated;
  Integer myTop;
  Integer Array myValues(1:Size);! but not hidden
                                     ! Methods;
                                      Operations encapsulated;
                                     ! Initialization;
  Procedure Init;
                                     ! 1-line methods need no 'end';
    myTop := 0;
  Boolean Procedure IsEmpty; ! Functions are typed procs; IsEmpty:= myTop < 1; ! Assign RV to function name
                                     ! Assign RV to function name;
  Boolean Procedure IsFull;
   IsFull:= myTop >= Size;
  ! ... continued on next page ... ;
```

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Example: An IntStack Class (ii)

```
IntStack.sim (continued)
 Procedure Push (Value);
                            ! Methods can have parameters;
                            ! Parameters precede 'Begin';
    Integer Value;
                            ! Methods with multiple statements;
  Begin
   myTop:= myTop + 1;
                            ! must be 'wrapped' in a block;
   myValues(myTop):= Value;
  End of Push;
 Procedure Pop;
  myTop:= myTop - 1;
 Integer Procedure Top;
   Top:= myValues(myTop);
                            ! Life: code following methods;
 Init;
                             is executed on object creation;
End of Stack;
Simula-67
                                    : nothing prevented code from:
```

- accessing the class's implementation details; or
- violating the intent of the creator of the class.



Simula: Using the IntStack Class

To use the class, we write something like this:

```
! IntStackTest.sim
Begin
 Ref(IntStack) S1, S2;
                                 ! Reference (pointer) variables;
  Integer value;
                                  ! Normal variable;
  S1 :- New IntStack(8);
                                  ! Special reference assignment;
  While Not S1. IsFull Do Begin
                                 ! Dot-notation for messages;
                                    no args, no parentheses;
   Value:= InInt;
                                  ! args require parentheses;
    S1.Push (Value);
  End;
  S2 :- S1;
                                  ! Reference assignment
  While Not S2. Is Empty Do Begin
    OutInt(S2.Top);
    S2.Pop;
  End;
End of Program;
```



Smalltalk

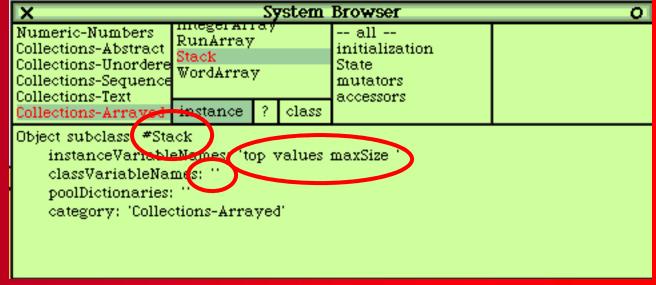
Smalltalk took the ideas of Simula and extended them:

Everything (including programs) are

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- Attribute variables are
- A large predefined *class library* is provided
- A GUI integrated development environment (IDE) is provided

Building a Smalltalk Stack class is simple, as most of the syntax is autogenerated by the IDE:





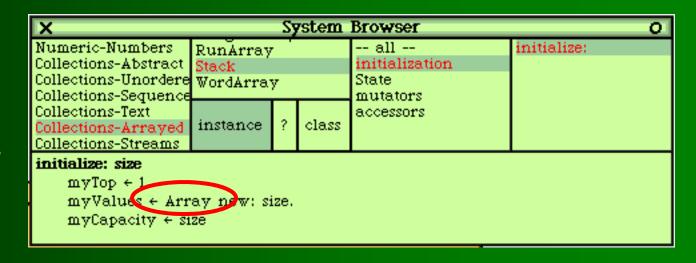
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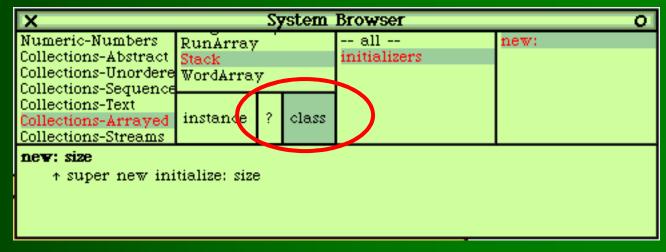
Smalltalk Operations

Adding *Stack*operations to our
class is also easy:
as the IDE makes
templates for us

Smalltalk *Arrays* store *Objects* (more later)...

new: is a message we send the class (i.e., a class method):

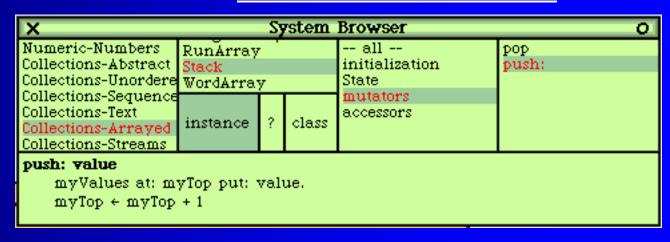




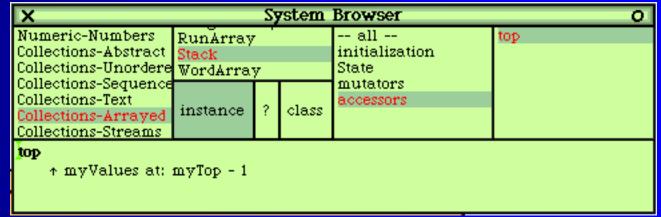


Smalltalk Operations (ii)

Mutators like *push* and *pop* are equally easy:



Other Stack
operations
(isEmpty, isFull,
top) are just as
easy...

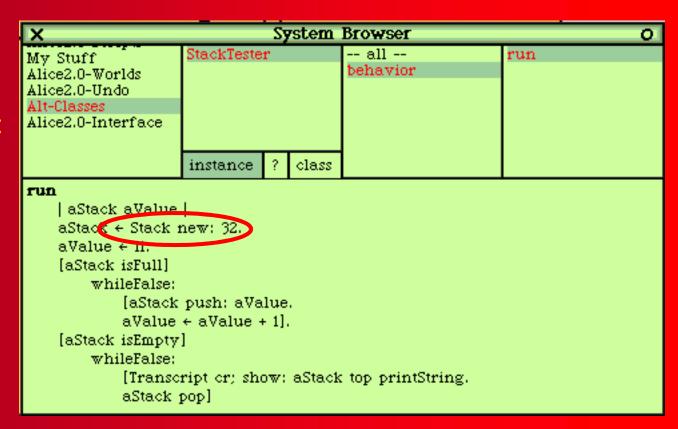


Smalltalk "Programs"

"Programs" are classes with ____ and ___ methods:

Note that we send new: to the class, not an instance of the class (i.e., an object)...

To run such programs:



X Workspace O
st := StackTester new run.



C++ Classes

In 1986,

```
// IntStack.h
class IntStack {
public:
   IntStack(int size);
   bool isEmpty() const;
   bool isFull() const;
   void push(int value);
   void pop();
   int top() const;
private:
   int myTop, myCapacity;
   int * myValues;
};
```

```
// IntStack.h (cont'd)
inline IntStack::IntStack(int size){
  myTop = -1;
  myCapacity = size;
  myValues = new int[size];
inline
bool IntStack::isEmpty() const{
  return myTop < 0;
inline void
IntStack::push(int value) const {
  myValues[++myTop] = value;
// ... other operation definitions ...
```

More complex operations can be separately compiled...



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Using C++ Classes

C++ objects can be

```
// IntStackTest1.cpp
#include "IntStack.h"
int main() {
 IntStack s(8);
 int aValue;
 while ( !s.isFull() ) {
  cin >> aValue;
  s.push(aValue);
 while ( !s.isEmpty() ) {
  cout << s.top() << ' ';
   s.pop();
```

```
// IntStackTest2.cpp
#include "IntStack.h"
int main() {
 IntStack* s = new IntStack(8);
 int aValue;
 while ( !s->isFull() ) {
   cin >> aValue;
   s->push(aValue);
 while ( !s->isEmpty() )
   cout << s->top() << ' ';
   s->pop();
```

n _____languages, objects *must* be dynamically allocated.

Java Classes

In 1993,

```
// IntStack.java
public class IntStack {
  public IntStack(int size) {
    myTop = -1;
    myCapacity = size;
    myValues = new int[size];
  }
  public boolean isEmpty() {
    return myTop < 0;
  }
  public boolean isFull() {
    return myTop >= myCapacity;
  }
```

, based on C++ and Smalltalk.

```
// IntStack.java (cont'd)
public void push(int value) {
   myValues[++myTop] = value;
public void pop() {
   --myTop;
public int top() {
   return myValues[myTop];
private int
                myTop,
                myCapacity;
private int [] myValues;
```

Java mixes



Using Java Classes

Like Smalltalk, Java objects must be dynamically allocated:

Class variables are

so dot notation is used (vs. C++ ->)

Also like Smalltalk, every Java subprog.

```
// IntStackTest.java
import IntStack;
class IntStackTest {
  public static void main(String [] args) {
  IntStack s = new IntStack(8);
  int aValue = 11;
  while ( !s.isFull() ) {
    s.push(aValue);
    aValue++;
  while ( !s.isEmpty() ) {
   System.out.println( s.top() );
   s.pop();
```

But most of Java's

is more similar to C++ than Smalltalk...



C++ Templates

C++ classes can have _

(→ Ada generic packages):

```
// Stack.h
template<class Item>
class Stack {
public:
   Stack(int size);
   bool isEmpty() const;
  bool isFull() const;
   void push(Item value);
   void pop();
   Item top() const;
private:
   int myTop, myCapacity;
   Item * myValues;
};
#include "Stack.tpp"
```

```
// Stack.tpp
template<class Item>
inline Stack<Item>::Stack(int size) {
  myTop = -1;
  myCapacity = size;
 myValues = new Item[size];
template<class Item>
inline bool Stack<Item>::isEmpty() {
  return myTop < 0;</pre>
template<class Item>
inline Item Stack<Item>::top() {
  return myValues[myTop];
// ... isFull, push, pop, ...
```

This is pretty clunky compared to Ada's generic mechanism...



C++ Template Instantiation

```
A template is a _
```

```
// StackTest1.cpp
#include "Stack.h"
int main() {
 Stack<int> s(32);
 int aValue;
 while ( !s.isFull() ) {
  cin >> aValue;
  s.push(aValue);
 while ( !s.isEmpty() ) {
   cout << s.top() << ' ';
   s.pop();
```

```
// StackTest1.cpp
#include "Stack.h"
#include "Student.h"
int main() {
  Stack<Student> s(32);
  Student aValue;
 while ( !s.isFull() ) {
   cin >> aValue;
   s.push(aValue);
 while ( !s.isEmpty() )
   cout << s.top() << ' ';
   s.pop();
```

Templates are esp. useful in creating classes that contain other objects.

Old Java Containers

Prior to Java 1.5, Java had no generics, but all classes have a common ancestor *Object*, so ______ were used:

```
// Stack.java
public class Stack {
 public Stack(int size) {
  myTop = -1;
  myCapacity = size;
  myValues = new Object[size];
 public boolean isEmpty() {
  return myTop < 0;</pre>
 public boolean isFull()
  return myTop >= myCapacity-1;
```

This can store any object



Old Java: Using Object Containers

Since our Stack stores Objects...

1. We can't store a primitive type (e.g., *int*) directly, but Java provides a

for each primitive type...

2. ____ must be used to retrieve the stored values...

```
// StackTest.java
import Stack;
class StackTest {
  public static void main(String [] args) {
  Stack s = new Stack(8);
  int aValue = 11;
  while (!s.isFull()) {
    s.push((new Integer(aValue)));
    aValue++;
  while (!s.isEmpty())
   Integer anInteger = (Integer))s.top();
   int anInt = anInteger.intvalue();
   System.out.println(anInt);
   s.pop();
```

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Newer Java Containers

Java 1.5 added and

```
// Stack.java
public class Stack<Item> {
 public Stack(int size) {
  myTop = -1;
  myCapacity = size;
  myValues = (Item[])
              new Object[size];
 public boolean isEmpty() {
  return myTop < 0;</pre>
 public boolean isFull() {
  return myTop >= myCapacity-1;
```

```
// Stack.java (cont'd)
public void push(Item value) {
  myValues[++myTop] = value;
public void pop() {
  --myTop;
public Item top() {
  return myValues[myTop];
 private int myTop,
                  myCapacity;
private Item [] myValues;
```

Such a *Stack* can still store



Java: Using Generic Containers

Since Stack<Item> stores Objects...

1. We can't pass a primitive type arg (e.g., *int*), but we can pass a _____.

2. ____ lets us pass primitive-type-values as args.

3. _____ lets us retrieve *Item*s as primitive type-vals.

```
// StackTester.java
import Stack;
class StackTester {
  public static void main(String [] args) {
 Stack<Integer> s = new Stack<>(8);
  int aValue = 11:
  while ( !s isFull() ) {
    s.push( aValue ); // int auto-boxed
    avaiue 11;
  while ( !s.isEmpty() ) {
  int anInt = s.top(); // auto-unboxed
   System.out.println(anInt);
   s.pop();
```

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Summary

To achieve the goals of abstract	data typing, we need:
: data and o	perations in;
- : the ability to	
Two different mechanisms have	evolved for doing so:
-The: a container for sto	oring a data-type and its operations
(aka the).	
-The: a type-constructor	that stores data and operations.
The class is better for	•
the module is better for	<u>. </u>
Containers are packages or clas	ses that
	support stronglytyped containers.
- Generic containers can be built in	most modern languages.
	A STATE OF THE STA