

# Control Structures

Programming Languages  
CS 214

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# Spaghetti Coding

In the early 1960s, \_\_\_\_\_ was common practice,  
whether using a HLL or a formal model like the RAM...

Example: What does this “spaghetti style” C function do?

```
double f(double n) {  
    double x = y = 1.0;  
    if (n < 2.0) goto label2;  
label1: if (x > n) goto label2;  
    y *= x;  
    x++;  
    goto label1;  
label2: return y;  
}
```

→ The \_\_\_\_\_

→ Such code was \_\_\_\_\_ *to maintain...*

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## Control Structures

In 1968, \_\_\_\_\_ published “Goto Considered Harmful”  
-- a letter suggested the *goto* should be outlawed because it encouraged undisciplined coding (the letter raised a furor).

Language designers began building \_\_\_\_\_  
-- statements whose syntax made control-flow obvious:

- |       |          |                |          |
|-------|----------|----------------|----------|
| •If   | Fortran  | •If-Then-Else  | COBOL    |
| •Case | Algol-W  | •If-Then-Elsif | Algol-68 |
| •For  | Algol-60 | •While         | Pascal   |
| •Do   | COBOL    |                |          |

With Pascal (1970), all of these were available in 1 language, resulting in a new coding style: \_\_\_\_\_.

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## Structured Programming

Structured Programming emphasized \_\_\_\_\_, through:

- Use of \_\_\_\_\_
- Use of \_\_\_\_\_
- Use of \_\_\_\_\_ (indentation, blank lines).

```
double factorial(double n)
{
    double result = 1.0;

    for (int count = 2; count <= n; count++)
        result *= count;

    return result;
}
```

With structured programming, \_\_\_\_\_!

The resulting programs were *less expensive to maintain*.

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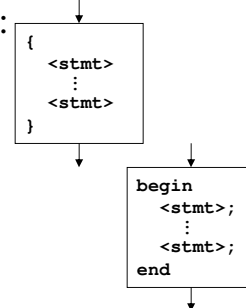
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## Sequential Execution

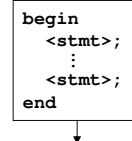
A C/C++ block has \_\_\_\_\_ statements:

$\langle \text{block-stmt} \rangle ::= \{ \langle \text{stmt-list} \rangle \}$   
 $\langle \text{stmt-list} \rangle ::= \langle \text{stmt} \rangle \langle \text{stmt-list} \rangle \mid \epsilon$



An Ada block has \_\_\_\_\_ stmts:

$\langle \text{block-stmt} \rangle ::= \text{begin } \langle \text{stmt-list} \rangle \text{ end}$   
 $\langle \text{stmt-list} \rangle ::= \langle \text{stmt} \rangle \langle \text{more-stmts} \rangle$   
 $\langle \text{more-stmts} \rangle ::= \langle \text{stmt} \rangle \langle \text{more-stmts} \rangle \mid \epsilon$



The block is the control structure for \_\_\_\_\_,  
the default control structure in imperative languages.

The guiding principle for control structures is:

\_\_\_\_\_.

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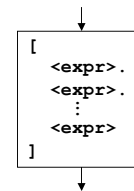
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## Smalltalk

Smalltalk also has a *block* construct, but it is an \_\_\_\_\_:

$\langle \text{block-object} \rangle ::= [ \langle \text{params} \rangle \langle \text{locals} \rangle \langle \text{expr-list} \rangle ]$   
 $\langle \text{params} \rangle ::= \langle \text{param-list} \rangle 'l' \mid \epsilon$   
 $\langle \text{param-list} \rangle ::= : \text{id} \langle \text{param-list} \rangle \mid \epsilon$   
 $\langle \text{locals} \rangle ::= 'l' \langle \text{id-list} \rangle 'l' \mid \epsilon$   
 $\langle \text{id-list} \rangle ::= \text{id} \langle \text{id-list} \rangle \mid \epsilon$   
 $\langle \text{expr-list} \rangle ::= \langle \text{expr} \rangle \langle \text{more-exprs} \rangle \mid \epsilon$   
 $\langle \text{more-exprs} \rangle ::= . \langle \text{expr} \rangle \langle \text{more-exprs} \rangle \mid \epsilon$



Smalltalk computations consist of *messages* sent to *objects*:

$[2 + 1] \text{ value} \rightarrow 3$

Like C/C++, a Smalltalk *block* can declare local variables;  
but as an object, a Smalltalk *block* can also have \_\_\_\_\_:

$[i \mid i + 1] \text{ value: } 2 \rightarrow 3$

$| \text{aBlock} | \text{aBlock} := [:x :y | (x*x) + (y*y)] . \text{aBlock value: } 3 \text{ value: } 4 \rightarrow 25$

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## Lisp

The expressions in the “body” of a Lisp function are executed sequentially, by default, with the value of the function being the value of \_\_\_\_\_:

```
(defun summation (n)
  (setq t1 (+ n 1))
  (setq t2 (* n t1))
  (setq t3 (/ t2 2)))

summation
(summation 100)
5050
```

Of course, *summation()* can be written more succinctly:

```
(defun summation (n)
  _____)
```

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## Lisp (ii)

Some Lisp function-arguments must be a single expression.

Lisp’s \_\_\_\_\_ function can be used to execute several expressions sequentially, much like other languages’ *block*:

```
(defun summation (n)
  (if (<= n 0)
      (progn
        (message "summation(n): n must be positive...")
        0)
      (/ (* (+ n 1) n) 2)))
```

The *progn* function returns the value of its \_\_\_\_\_.

Lisp also has sequential *prog1* and *prog2* functions, that return the values of the 1<sup>st</sup> and 2<sup>nd</sup> expressions, respectively.

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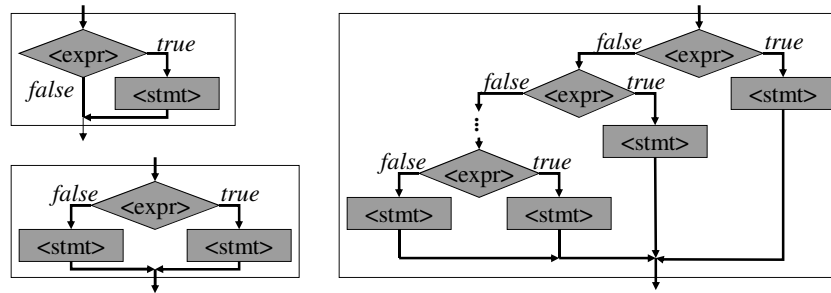
## Selective Execution

... lets us \_\_\_\_\_.

The *If Statement* provides selective execution:

```
<if-statement> ::= if ( <expr> ) <stmt> <else-part>
<else-part>    ::= else <stmt> | ε
```

These rules permit three different forms of flow control:



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## Examples

These three forms allow us to use selective execution in whatever manner is appropriate to solve a given problem:

\_\_\_\_\_ Logic:

```
if (numValues != 0)
    avg = sum / numValues;
```

\_\_\_\_\_ Logic:

```
if (first < second)
    min = first;
else
    min = second;
```

\_\_\_\_\_ Logic:

```
if (score > 89)
    grade = 'A';
else if (score > 79)
    grade = 'B';
else if (score > 69)
    grade = 'C';
else if (score > 59)
    grade = 'D';
else
    grade = 'F';
```

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## The Dangling Else Problem

Every language designer must resolve the question of how to associate a “dangling else” following nested if statements...

```

if Condition1 then
  if Condition2 then
    Statement1
  else
    Statement2

```

?

The problem occurs in languages with \_\_\_\_\_.

→ Such a statement can be \_\_\_\_\_ in two different ways.

There are two different approaches to resolving the question:

- Add a semantic rule to resolve the ambiguity; vs.
- Design a statement whose syntax is not ambiguous.

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## Using Semantics

Languages from the 1970s (Pascal, C) tended to use simple but ambiguous grammars:

```

<if-stmt>      ::= if ( <expr> ) <stmt> <else-part>
<else-part>   ::= else <stmt> | ε

```

plus a semantic rule:

```

if ( Condition1 )
  if ( Condition2 )
    Statement1
  else
    Statement2

```

if ( Condition<sub>1</sub> )  
 {  
     if ( Condition<sub>2</sub> )  
         Statement<sub>1</sub>  
 }  
 else  
     Statement<sub>2</sub>

Block statements provided a way to circumvent the rule. Newer C-family languages (C++, Java) have inherited this.

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## Using Syntax

Newer languages tend to use \_\_\_\_\_:

```
<if-stmt> ::= if ( <expr> ) <stmt-list> <else-part> end if
<else-part> ::= else <stmt-list> | ε
<stmt-list> ::= <stmt> <stmt-list> | ε
```

Terminating an *if* with an *end if* “closes” the most recent *else*, eliminating the ambiguity without any semantic rules:

<pre>if ( Condition<sub>1</sub> )   if ( Condition<sub>2</sub> )     StmtList<sub>1</sub>   else     StmtList<sub>2</sub>   end if end if</pre>	<pre>if ( Condition<sub>1</sub> )   if ( Condition<sub>2</sub> )     StmtList<sub>1</sub>   end if else   StmtList<sub>2</sub> end if</pre>
---	---

Ada, Fortran, Modula-2, ... use this approach.

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## Using Syntax (ii)

Perl uses a (different) syntax solution:

```
<if-stmt> ::= if ( <expr> ) <block> <else-part>
<else-part> ::= else <block> | ε
<block> ::= { <stmt-list> }
```

By requiring each branch of an *if* to be a *block*, \_\_\_\_\_, eliminating the ambiguity:

<pre>if ( Condition<sub>1</sub> ) {   if ( Condition<sub>2</sub> ) {     StmtList<sub>1</sub>   } else {     StmtList<sub>2</sub>   } }</pre>	<pre>if ( Condition<sub>1</sub> ) {   if ( Condition<sub>2</sub> ) {     StmtList<sub>1</sub>   } } else {   StmtList<sub>2</sub> }</pre>
---	---

The end of the block serves to terminate the nested *if*.

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## Aesthetics

Multibranch selection can get clumsy using *end if*:

```
if ( Condition1 )  
    StmtList1  
else if ( Condition2 )  
    StmtList2  
    else if ( Condition3 )  
        StmtList3  
        else  
            StmtList4  
        end if  
    end if  
end if
```

```
if ( Condition1 )  
    StmtList1  
elseif ( Condition2 )  
    StmtList2  
elseif ( Condition3 )  
    StmtList3  
else  
    StmtList4  
end if
```

To avoid this problem, Algol-68 added the *elif* keyword that, substituted for *else if*, \_\_\_\_\_.

Modula-2 and Ada replaced the error-prone *elif* with \_\_\_\_\_.

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## Exercise

Write a BNF for Ada *if*-statements. Sample statements:

```
if numValues <> 0 then  
    avg := sum / numValues;  
end if;
```

```
if first < second then  
    min := first;  
    max := second;  
else  
    min := second;  
    max := first;  
end if;
```

```
if score > 89 then  
    grade := 'A';  
elseif score > 79 then  
    grade := 'B';  
elseif score > 69 then  
    grade := 'C';  
elseif score > 59 then  
    grade := 'D';  
else  
    grade := 'F';  
end if;
```

<Ada-if-stmt> ::= \_\_\_\_\_  
\_\_\_\_\_ ::= \_\_\_\_\_  
\_\_\_\_\_ ::= \_\_\_\_\_

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## Lisp's *if*

Lisp provides an \_\_\_\_\_ as one of its expressions:

```
<if-expr>      ::= ( if <predicate> <expr> <opt-expr> )  
<opt-expr>    ::= <expr> | ε
```

Semantics: If the <predicate> evaluates to non-nil (i.e., not ()),  
the <expr> is evaluated and its value returned;  
else the <opt-expr> is evaluated and its value returned.

```
(if (> score 89)  
    (setq grade "A")  
    (if (> score 79)  
        (setq grade "B")  
        (if (> score 69)  
            (setq grade "C")  
            (if (> score 59)  
                (setq grade "D")  
                (setq grade "F")))))
```

It is not unusual for a Lisp  
expression to end with ))))

\_\_\_\_\_

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## Selection in Smalltalk

Smalltalk provides various \_\_\_\_\_  
that can be sent to \_\_\_\_\_...

```
<selection-msg> ::= <ifT-msg> | <ifF-msg> | <ifTF-msg> | <ifFT-msg>  
<ifT-msg>      ::= ifTrue: <block>  
<ifF-msg>      ::= ifFalse: <block>  
<ifTF-msg>     ::= ifTrue: <block> ifFalse: <block>  
<ifFT-msg>     ::= ifFalse: <block> ifTrue: <block>
```

```
n ~= 0  
ifTrue: [ avg := sum / n ]
```

```
first < second  
ifTrue: [ min := first]  
ifFalse: [ min := second]
```

These four are the only selection  
messages Smalltalk provides.

```
score > 89  
ifTrue: [grade:= 'A']  
ifFalse: [  
    score > 79  
    ifTrue: [grade:= 'B']  
    ifFalse: [  
        score > 69  
        ifTrue: [grade:= 'C']  
        ifFalse: [ ... ] ] ]
```

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## Problem: Non-Uniform Execution

```
if (score > 89)      ← ____ comparison to get here
    grade = 'A';
else if (score > 79) ← ____ comparisons to get here
    grade = 'B';
else if (score > 69) ← ____ comparisons to get here
    grade = 'C';
else if (score > 59) ← ____ comparisons to get here...
    grade = 'D';
else                ← ... and here
    grade = 'F';
```

The times to execute different branches are \_\_\_\_\_:

- The 1<sup>st</sup> <stmt> executes after \_\_\_\_ comparison.
- The n<sup>th</sup> and final <stmt> execute after \_\_\_\_ comparisons.

The time to execute successive branches increases \_\_\_\_\_.

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## The Switch Statement

The switch statement provides \_\_\_\_\_.

```
<switch>      ::= switch ( <expr> ) { <pair-list> <opt-default> }
<pair-list>   ::= <case-list> <stmt-list> <pair-list> | ε
<case-list>   ::= case <literal> : <case-list> | ε
<opt-default> ::= default: <stmt-list> | ε
```

Rewriting our grade program:

Note: If you neglect to supply *break* statements, control by default flows \_\_\_\_\_ through the *switch* statement.

The *break* is a \_\_\_\_\_ statement...

```
switch (score / 10) {
    case 9: case 10:
        grade = 'A'; break;
    case 8:
        grade = 'B'; break;
    case 7:
        grade = 'C'; break;
    case 6:
        grade = 'D'; break;
    default:
        grade = 'F';
}
```

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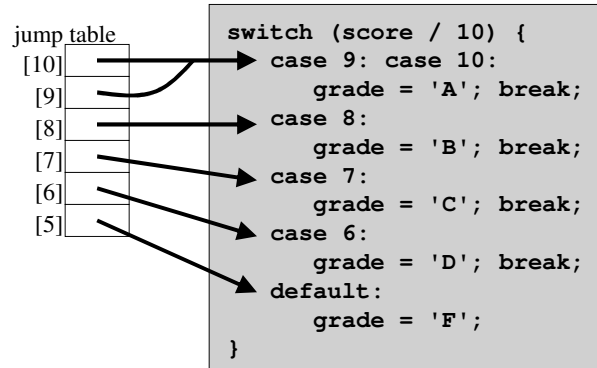
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## Uniform Execution Time

Compiled *switch/case* statements achieve uniform response time via a \_\_\_\_\_, that stores the address of each branch.



This is simplified a bit, but it gives the general idea...

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## Uniform Execution Time (ii)

With a jump table, a compiler can translate a *switch/case* to something like this:

```
// code to evaluate <expr>
// and store it in register R

    cmp R, #highLiteral
    jle lowerTest
    mov #lowLiteral-1, R
    jmp makeTheJump

lowerTest:
    cmp R, #lowLiteral
    jge makeTheJump
    mov #lowLiteral-1, R

makeTheJump:
    mov jumpTable[R], PC

// branches of the switch
```

For non-default branches, a *switch/case* needs \_\_\_\_\_ and \_\_\_\_\_ to find the branch.

When a multibranch *if* does \_\_\_\_\_, a *switch* is probably faster.

A compiler spends \_\_\_\_\_ time and space (to build the jump table) to decrease the average \_\_\_\_\_ time needed to find a branch.

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## The Case Statement

The *switch* is a descendent of the \_\_\_\_\_ statement (Algol-W).  
Only C-family languages use the *switch* syntax.

Unlike the *switch*, a *case* statement \_\_\_\_\_  
\_\_\_\_\_ behavior.

Most *case* stmts also let you use literal \_\_\_\_\_ and \_\_\_\_\_:

Ada uses the *when* keyword to begin each <literal-list>, and uses the => symbol to terminate each literal-list.

```
case score / 10 of
  when 9, 10 =>
    grade = 'A';
  when 8 =>
    grade = 'B';
  when 7 =>
    grade = 'C';
  when 6 =>
    grade = 'D';
  when 0..5 =>
    grade = 'F';
  when others =>
    put_line("error...");
end case;
```

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## Exercise

Build a BNF for Ada's *case* statement.

- There must be at least one branch in the statement.
- A branch must contain at least one statement.
- The *when others* branch is optional, but must appear last.

<Ada-case> ::= \_\_\_\_\_

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## Lisp

Lisp provides a \_\_\_\_\_ that looks similar to a *case*.

```
<cond-expr> ::= ( cond <expr-pairs> )
<expr-pairs> ::= ( <predicate> <expr> ) <expr-pairs> | ε
```

However Lisp's *cond* uses  
arbitrary predicates  
(relational expressions)  
instead of literals.

```
(cond
  (> score 89) "A"
  (> score 79) "B"
  (> score 69) "C"
  (> score 59) "D"
  (t "F")
)
```

→ As a result, Lisp's *cond* cannot employ a jump table,  
so it has the same non-uniform execution time as an *if*.

The predicates are evaluated \_\_\_\_\_ until a true  
<predicate> is found; its <expr> is then evaluated.

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## Repetition

A third control structure is \_\_\_\_\_, or \_\_\_\_\_.

The C++ *while* loop is a \_\_\_\_\_ loop:

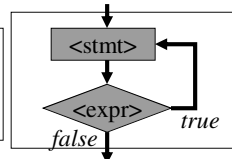
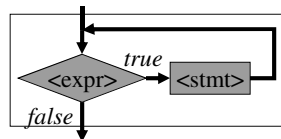
```
<while-stmt> ::= while ( <expr> ) <stmt>
```

but the *do* loop is a \_\_\_\_\_ loop:

```
<do-stmt> ::= do <stmt> while ( <expr> );
```

Which is which?

```
while ( <expr> )
  <stmt>
```



```
do
  <stmt>
while ( <expr> );
```

A pretest loop's <stmt> is executed \_\_\_\_\_ times (\_\_\_\_\_).

A posttest loop's <stmt> is executed \_\_\_\_\_ times (\_\_\_\_\_).

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## Counting Loops

Like most languages, C++ provides a \_\_\_\_ loop for counting:

```
<for-stmt> ::= for ( <opt-expr> ; <opt-expr> ; <opt-expr> ) <stmt>
```

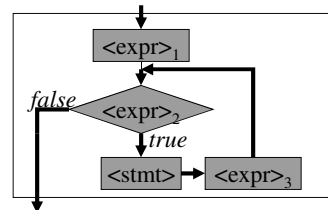
This provides unusual flexibility for an imperative language:

```
for (int i = 0; i <= 100; i++)
    <stmt> // do <stmt> 101 times; i = _____

for (double d = -0.5; d <= 0.5; d += 0.1)
    <stmt> // do <stmt> 11 times; d = _____

for (Node * ptr = myHead; ptr != myTail; ptr = ptr->next)
    <stmt> // do <stmt> _____
```

In most languages,  
the counting loop is  
a \_\_\_\_\_ loop:



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## Unrestricted Loops

Most modern languages also support an \_\_\_\_\_.

–Such loops have \_\_\_\_\_.

–All of the C/C++ loops can be made to behave this way.

```
for (;;)      while (true)    do
    <stmt>      <stmt>        <stmt>
                        while (true);
```

–The language usually provides a statement to exit such loops.

–Unrestricted loops can be structured as \_\_\_\_\_,  
\_\_\_\_\_, or \_\_\_\_\_ loops:

```
for (;;) {
    if (<expr>) break;
    <stmt>2
}
```

```
for (;;) {
    <stmt>1
    if (<expr>) break;
}
```

```
for (;;) {
    <stmt>1
    if (<expr>) break;
    <stmt>2
}
```

– <stmt>1 executes \_\_\_\_ times; <stmt>2 executes \_\_\_\_ times...

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## Ada

Ada provides \_\_\_\_\_, \_\_\_\_\_, and \_\_\_\_\_ loops:

```
for i in 1..100 loop
  ...
end loop;
```

```
for i in reverse 1..100 loop
  ...
end loop;
```

```
while i <= 100 loop
  ...
  i:= i+1;
end loop;
```

```
loop
  exit when i > 100;
  ...
  i:= i+1;
end loop;
```

Exercise: How would you build a BNF for Ada's loops?

<Ada-loop-stmt> ::= \_\_\_\_\_  
 \_\_\_\_\_ ::= \_\_\_\_\_  
 \_\_\_\_\_ ::= \_\_\_\_\_  
 \_\_\_\_\_ ::= \_\_\_\_\_  
 \_\_\_\_\_ ::= \_\_\_\_\_

What if you need a post-test loop, or to count by i != 1?

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## Smalltalk

Smalltalk provides \_\_\_\_\_:

```
<loop-expr> ::= <while-expr> | <times-expr> | <to-expr>
<while-expr> ::= <block> <while-msg> <block>
<while-msg> ::= whileTrue: | whileFalse:
<times-expr> ::= <intExpr> timesRepeat: <block>
<to-expr> ::= <numExpr> to: <numExpr> <opt-by> do: <block>
<opt-by> ::= by: <numExpr> | ε
```

```
[i <= 100] whileTrue:
[
  ...
  i:= i+1
]
```

```
[i > 100] whileFalse:
[
  ...
  i:= i+1
]
```

```
100 timesRepeat:
[
  ...
]
```

```
0 to: 100 do:
[
  ...
]
```

```
-0.5 to: 0.5 by: 0.1 do:
[
  ...
]
```

Under what circumstances should a given loop be used?

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## Lisp

Lisp has no loop functions, because anything that can be done by repetition can also be done using \_\_\_\_\_.

```
(defun f(n)
  ...
  (f(+ n 1))
)
```

```
(defun factorial(n)
  (if (< n 2)
      1
      (* n (factorial (- n 1)))))
)
```

Recursive functions can provide test-at-the-top, test-at-the-bottom, and test-in-the-middle behavior simply by varying

```
(defun f(n)
  (if (< n max)
      (f(+ n 1))
      <expr-list>)
)
```

```
(defun g(n)
  <expr-list>
  (if (< n max)
      (f(+ n 1)))
)
```

```
(defun h(n)
  <expr-list>
  (if (< n max)
      (f(+ n 1))
      <expr-list>)
)
```

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## Summary

There are three basic control structures:

- \_\_\_\_\_
- \_\_\_\_\_
- \_\_\_\_\_

Different kinds of languages accomplish these differently:

- Sequence is the default mode of control provided by the \_\_\_\_\_ construct of most languages (\_\_\_\_\_ in Lisp).
- Selection is accomplished via:
  - \_\_\_\_\_ (e.g., *if*, *switch* or *case*) controlled by boolean expressions in imperative languages
  - \_\_\_\_\_ (e.g., *if* and *cond* in Lisp) with boolean arguments in functional languages
  - \_\_\_\_\_ (*ifTrue:*, *ifFalse:*, ... in Smalltalk) sent to boolean objects in pure OO languages

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## Summary (ii)

- Repetition is accomplished via:

- \_\_\_\_\_ (e.g., *while*, *do*, *for*) controlled by boolean expressions in imperative languages
- \_\_\_\_\_ in functional languages
- \_\_\_\_\_ (*whileTrue:*, *timesRepeat:*, *to:by:do:*, ... in Smalltalk) sent to boolean (or numeric) objects in pure OO languages

These \_\_\_\_\_ are all we need to compute anything that can be computed (i.e., by a Turing machine).

Most of the other language constructs simply make the task of programming such computations \_\_\_\_\_.

