



SCHOOL OF
COMPUTING

Vishal M.D.

CH.SC.U4CSE2415

0

Week - 3

Date - 03/01/2026

Design and Analysis of Algorithm(23CSE211)

1. Merge Sort

Code:

```
#include <stdio.h>
#include <stdlib.h>

void merge(int arr[], int left, int mid, int right) {
    int i = left, j = mid + 1, k = 0;
    int *temp = (int *)malloc((right - left + 1) * sizeof(int));
    while (i <= mid && j <= right) {
        if (arr[i] <= arr[j]) {
            temp[k++] = arr[i++];
        } else {
            temp[k++] = arr[j++];
        }
    }
    while (i <= mid) {
        temp[k++] = arr[i++];
    }
    while (j <= right) {
        temp[k++] = arr[j++];
    }
    for (i = left, k = 0; i <= right; i++, k++) {
        arr[i] = temp[k];
    }
    free(temp);
}

void mergeSort(int arr[], int left, int right) {
    if (left < right) {
        int mid = left + (right - left) / 2;
        mergeSort(arr, left, mid);
        mergeSort(arr, mid + 1, right);
        merge(arr, left, mid, right);
    }
}

void printArray(int arr[], int n) {
    for (int i = 0; i < n; i++) {
        printf("%d ", arr[i]);
    }
    printf("\n");
}

int main() {
    int arr[] = {64, 34, 25, 12, 22, 11, 90};
    int n = sizeof(arr) / sizeof(arr[0]);
    printf("Original array: ");
    printArray(arr, n);
    mergeSort(arr, 0, n - 1);
    printf("Sorted array: ");
```

```
printArray(arr, n);
return 0;
}
```

Output:

```
● (base) → DAA gcc ./mergesort.c -o m
● (base) → DAA ./m
Original array: 64 34 25 12 22 11 90
Sorted array: 11 12 22 25 34 64 90
❖ (base) → DAA █
```

Space Complexity: $O(n \log n)$

Justification: The algorithm recursively divides the array into halves, which takes $\log n$ levels of division. At each level, it performs a linear merge process requiring n operations, resulting in a consistent $n \times \log n$ performance.

Time Complexity: $O(n \log n)$

Justification: Merge Sort requires an auxiliary array to temporarily hold elements during the merge process. This additional memory is proportional to the size of the input array.

2. Quick Sort

Code:

```
#include <stdio.h>

void swap(int *a, int *b)
{
    int temp = *a;
    *a = *b;
    *b = temp;
}

int partition(int arr[], int low, int high)
{
    int pivot = arr[high];
    int i = low - 1;

    for (int j = low; j < high; j++)
    {
        if (arr[j] < pivot)
        {
            i++;
            swap(&arr[i], &arr[j]);
        }
    }
    swap(&arr[i + 1], &arr[high]);
    return i + 1;
}

void quickSort(int arr[], int low, int high)
{
    if (low < high)
    {
        int pi = partition(arr, low, high);
        quickSort(arr, low, pi - 1);
        quickSort(arr, pi + 1, high);
    }
}

void printArray(int arr[], int n)
{
    for (int i = 0; i < n; i++)
        printf("%d ", arr[i]);
    printf("\n");
}

int main()
{
    int arr[] = {64, 34, 25, 12, 22, 11, 90};
    int n = sizeof(arr) / sizeof(arr[0]);
```

```
printf("Original array: ");  
printArray(arr, n);
```

```
quickSort(arr, 0, n - 1);
```

```
printf("Sorted array: ");  
printArray(arr, n);
```

```
return 0;  
}
```

Output:

```
● (base) → DAA ./q
Original array: 64 34 25 12 22 11 90
Sorted array: 11 12 22 25 34 64 90
◆ (base) → DAA █
```

Space Complexity: $O(n^2)$

Justification: This occurs when the pivot consistently picks the smallest or largest element (e.g., on a sorted array). This results in highly unbalanced partitions where one side has 0 elements and the other has $n-1$, leading to n recursive levels with $O(n)$ work each.

Time Complexity: $O(n)$

Justification: In the worst-case scenario of an unbalanced partition, the recursion stack depth increases from the ideal $\log n$ to n . Each recursive call stays on the stack until it finishes, requiring memory proportional to the number of elements.