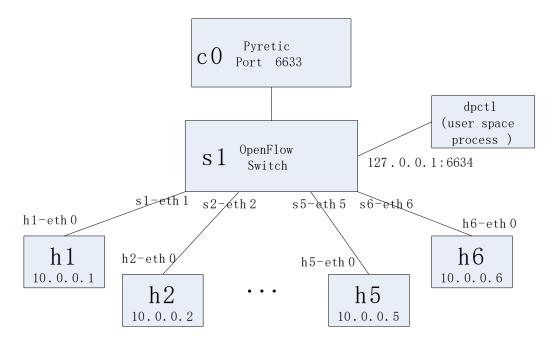
高级计算机网络第三次实验报告

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1. 实验目标

使用 Pyretic 实现第二层 Mac 地址上的防火墙。

2. 网络的拓扑结构



由上图所示的拓扑结构可以看出,本次实验的网络拓扑非常简单,6 个主机直接连接到交换机上面,控制器为 Pyretic 控制器。

3. 防火墙代码及相关的解释

#本策略使用的黑名单方法,黑名单中指定禁止的策略,其他的默认通过。

#导入相关的包以及相关模块

from pyretic.lib.corelib import *

from pyretic.lib.std import *

from pyretic.examples.pyretic_switch import act_like_switch

import os, csv

```
policy_file = "%s/pyretic/pyretic/examples/firewall-policies.csv" % os.environ[ 'HOME' ]
def main():
   #初始化
 not_allowed = none
 #从 csv 读取策略文件
 with open(policy_file, "r") as policy_content:
   dictReader = csv.DictReader(policy_content)
     #添加禁止策略,双向通信都禁止
   for d in dictReader:
   not_allowed = not_allowed +
((match(srcmac=MAC(d['mac_0']))&match(dstmac=MAC(d['mac_1'])))+
( match(srcmac=MAC(d['mac_1']))&match(dstmac=MAC(d['mac_0']))))
 #添加许可策略,表示为禁止策略取反。
 allowed = ~not_allowed
 #先输出一下当前规则
 print allowed
 #将允许的路由输入给作为 pyretic_switch 中的 act_like_switch
 return allowed >> act_like_switch()
```

4. 防火墙规则

```
id,mac_0,mac_1
1,00:00:00:00:00:01,00:00:00:00:04
```

```
2,00:00:00:00:02,00:00:00:00:05
3,00:00:00:00:03,00:00:00:00:06
```

如上表所示,为 firewall-policies 的规则,首列为规则编号,后面跟着的两个为不允许通讯的两个主机的 Mac 地址。

5. 实验及相关结果

5.1 启动控制器

将上述两个文件保存到 Pyretic/examples 的目录下之后,执行以下命令启动控制器。

pyretic.py -v high pyretic.examples.pyretic_firewall

执行之后可以观察到控制器已经启动(如图 5.1),没有报错说明防火墙程序没有问题。

5.2 创建拓扑

执行以下命令创建一个 6个节点连接到一个交换机上面的网络拓扑。

```
sudo mn --topo single,6 --mac --switch ovsk --controller remote
```

执行完成之后观察到如图 5.2 所示的界面,表示拓扑创建成功。

```
mininet@mininet ~/pyretic/pyretic/examples

*** Removing excess kernel datapaths

ps ax | egrep -o 'dp[0-9]+' | sed 's/dp/nl:/'

*** Removing OVS datapathsovs-vsctl list-br

*** Removing 3ll links of the pattern foo-ethX

ip link show | egrep -o '(\w+-eth\w+)'

*** Cleanup complete.

mininet@mininet:-/pyretic/pyretic/examples$ sudo mn --topo single,6 --mac --switch ovsk --controller remote

*** Creating network

*** Adding controller

*** Adding hosts:

hl h2 h3 h4 h5 h6

*** Adding switches:

sl

*** Adding links:

(h1, sl) (h2, sl) (h3, sl) (h4, sl) (h5, sl) (h6, sl)

*** Configuring hosts

hl h2 h3 h4 h5 h6

*** Starting controller

*** Starting outroller

*** Starting l switches

sl

*** Starting l switches

sl

*** Starting l switches
```

5.3 Pingall 测试

拓扑创建完成之后,执行 pingall 命令,观察到如下图 5,3 所示的界面,可以很明显的看出定义的防火墙策略已经生效。

5.4 单个交换机 Ping 测试

```
mininet> h1 ping h2
PING 10.0.0.2 (10.0.0.2) 56(84) bytes of data.
From 10.0.0.1 icmp_seq=1 Destination Host Unreachable
From 10.0.0.1 icmp_seq=2 Destination Host Unreachable
From 10.0.0.1 icmp_seq=3 Destination Host Unreachable
From 10.0.0.1 icmp_seq=4 Destination Host Unreachable
From 10.0.0.1 icmp_seq=5 Destination Host Unreachable
From 10.0.0.1 icmp_seq=6 Destination Host Unreachable
From 10.0.0.1 icmp_seq=6 Destination Host Unreachable
From 10.0.0.2 ping statistics ---
9 packets transmitted, 0 received, +6 errors, 100% packet loss, time 8046ms
pipe 3
mininet> h1 ping h3
PING 10.0.0.3 (10.0.0.3) 56(84) bytes of data.
64 bytes from 10.0.0.3: icmp_req=1 ttl=64 time=227 ms
64 bytes from 10.0.0.3: icmp_req=2 ttl=64 time=0.049 ms
64 bytes from 10.0.0.3: icmp_req=3 ttl=64 time=0.048 ms
64 bytes from 10.0.0.3: icmp_req=4 ttl=64 time=0.152 ms
^C
--- 10.0.0.3 ping statistics ---
4 packets transmitted, 4 received, 0% packet loss, time 3000ms
rtt min/avg/max/mdev = 0.048/56.860/227.191/98.340 ms
mininet> h1 ping h5
PING 10.0.0.5 (10.0.0.5) 56(84) bytes of data.
64 bytes from 10.0.0.5: icmp_req=2 ttl=64 time=152 ms
64 bytes from 10.0.0.5: icmp_req=2 ttl=64 time=0.051 ms
64 bytes from 10.0.0.5: icmp_req=2 ttl=64 time=0.051 ms
64 bytes from 10.0.0.5: icmp_req=2 ttl=64 time=0.051 ms
64 bytes from 10.0.0.5: icmp_req=3 ttl=64 time=0.054 ms
7C
--- 10.0.0.5 ping statistics ---
3 packets transmitted, 3 received, 0% packet loss, time 2000ms
rtt min/avg/max/mdev = 0.046/50.823/152.372/71.805 ms
mininet>
```

可以看出, h1 无法与 h2\h4\h6 通信, 可以 ping 通 h3\h5

6. 总结

本次实验使用 Pyretic 控制器,使用控制器实现了 Mac 层上的防火墙相关的策略。也很容易从本次 Project 中看出,使用 SDN 网络实现防火墙策略的管理比传统网络更加容易。