

STL: C++ : part 2:

lists

list<int> ls;

ls.push_back(2); // {2}

ls.emplace_back(4); // {2, 4}

cheap ← ls.push_front(5); // {5, 2, 4}; (Insert in a vector is costlier)

ls.emplace_front(1); // {1, 5, 2, 4}

// rest begin, end, begin, end, clear, insert, size, swap
are same as vectors

Deque

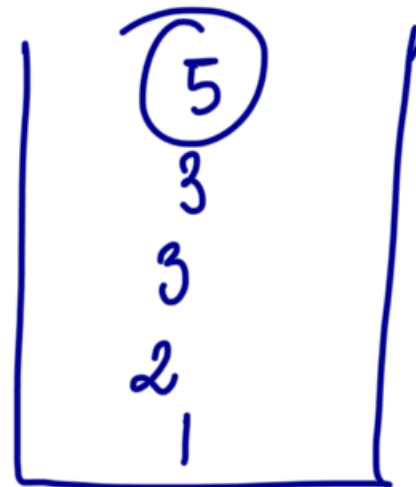
deque<int> dq;

`dq.push-back(1);` $\rightarrow \{1\}$
`dq.emplace-back(2);` $\rightarrow \{1, 2\}$
`dq.push-front(4);` $\rightarrow \{4, 1, 2\}$
`dq.emplace-front(3);` $\rightarrow \{3, 4, 1, 2\}$
`dq.pop-back();` $\rightarrow \{3, 4, 1\}$
`dq.pop-front();` $\rightarrow \{4, 1\}$
`dq.back();` $\rightarrow 1$
`dq.front();` $\rightarrow 4$

Just: `begin`, `end`, `begin`,
`rend`, `clear`, `insert`, `size`,
`swap` are same as
 vectors

Stack: Last in first out (LIFO)

`stack<int> st;`
`st.push(1);`
`st.push(2);`
`st.push(3);`
`st.push(3);`
`st.emplace(5);`



indexing not allowed

`cout << st.top();` $\rightarrow 5$

3 functions: `push`, `pop`, `top` $\rightarrow O(1)$

all things happen

st.pop(); // 5 is deleted in constant time

st.top(); $\rightarrow 3$

cout << st.size(); $\rightarrow 4$

cout << st.empty(); \rightarrow false

stack <int> st1, st2;
st1.swap(st2);

Queue : Fifo : first in first out

queue <int> q;

q.push(1); q.push(2); q.emplace(4);

q.back() = 5;

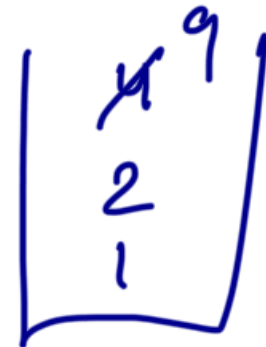
cout << q.back(); $\rightarrow 9$

cout << q.front(); $\rightarrow 1$

q.pop \rightarrow

9
2

// size, swap, empty same as stack



cout << q.front(); → 2

|| all operations in constant time

priority Queue:

priority_queue <int> pq;

pq.push(5);

pq.push(2);

pq.push(8);

pq.push(10); {10, 8, 5, 2}

(max heap)

10
8
5
2

tree data structure,
not linear.

cout << pq.top(); → 10

pq.pop(); →

8
5
2

push, top, pop: main

size, swap, empty: side func.

|| Minimum heap → minimum priority queue

priority_que <int, vector<int>, greater<int>> pq;

pq.push(5);

pq.push(2);

2
5

pq.push(8);
pq.emplace(10);



push, pop $\rightarrow O(\log n)$
top $\rightarrow O(1)$

cout << pq.top(); $\rightarrow 2$

// Set: Sorted order, unique

$O(\log N)$

set<int> st;

st.insert(1);

st.emplace(2);

st.insert(2);

st.insert(4);

st.insert(3);

{1}

{1, 2}

{1, 2}

{1, 2, 4}

{1, 2, 3, 4}



tree, not a linear container

// {1, 2, 3, 4, 5}

(begin(), end(), rbegin(), rend(), size())

auto it = st.find(3) \rightarrow returns an iterator that points to 3

auto it = st.find(6) \rightarrow returns st.end()

st.erase(5)

// erases 5, logarithmic time

int cnt = st.count(1); $\nearrow 1$
 $\searrow 0$

auto it = st.find(3);
st.erase(it);

// constant time

// {1, 2, 3, 4, 5}

auto it1 = st.find(2);

auto it2 = st.find(4);

st.erase(it1, it2);

[first, last)

(size, empty,
swap, begin: vector)

// lower_bound() & upper_bound() work in the same way
as in vector it does.

auto it = st.lower_bound(2);
auto it = st.upper_bound(3);

(do more after
learning binary
search)
