The M Specification

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1 Syntax

```
 \langle character \rangle \quad ::= \mathbb{C} \supseteq \langle special + newline + whitespace \rangle 
 \langle special \rangle \quad ::= \{';', ''', '(', ')'\} 
 \langle comment \rangle \quad ::= ';' \langle character - newline \rangle^* \langle newline \rangle 
 \langle symbol \rangle \quad ::= \langle character - special - whitespace \rangle^* 
 \mid ''' (\langle character - ''' \rangle \mid ' \setminus ' \langle character \rangle)^* ''' 
 \langle expression \rangle \quad ::= \langle symbol \rangle 
 \mid '(' \langle expression \rangle^* ')' 
 \langle program \rangle \quad ::= \langle expression \rangle^*
```

Figure 1: The M grammar in EBNF.

1.1 Characters

M can use any character set which encodes the four special characters, a backslash, a newline character, and a whitespace character. For consistency, all M code samples in this specification use the ASCII character set.

1.2 Specials

The four special characters are reserved for use in other productions. This is done to ensure that symbols such as this() do not include any special characters which are meant be used for other syntactic constructs. Note that the backslash is not included in the list of special characters as it is only used inside of a literal symbol.

1.3 Comments

Comments in M begin with a semicolon and last until the end of the line. They are ignored and discarded like whitespace and newlines, as they are only intended to be used for explaining code.

1.4 Symbols

Symbols in M have two forms, inline and literal. Inline symbols are strings terminated by whitespace or a special symbol, and should be used by default. Literal symbols are escapable strings surrounded by quotes, and should only be used when a symbol is impossible to represent using inline symbols (for example, the symbol "()").

1.5 Expressions

Expressions in M have two forms, symbols and lists. Symbol expressions are identical to symbols defined in section 1.4. Lists expressions are simply lists of expressions surrounded by matching parentheses.

1.6 Programs

M programs are lists of expressions. It is unspecified how these expressions are stored, so they can be in memory, in a single file, in multiple files, on the internet, or any other method which is convenient.

2 Semantics

$$(Function) \frac{1}{(\mathbf{fn} \times e) \Downarrow \lambda x.e} \qquad \frac{\langle f, \Gamma \rangle \Downarrow \lambda x.e \qquad \langle a, \Gamma \rangle \Downarrow i \qquad \langle e[x/i] \rangle \Downarrow v}{\langle (\mathbf{f} \text{ a}), \Gamma \rangle \Downarrow v} \qquad (Apply-Function) \\ (Impure) \frac{\langle e, \Gamma \rangle \Downarrow v}{\langle (\mathbf{impure} \text{ e}), \Gamma \rangle \Downarrow v} \qquad \frac{\langle f, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ a}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ (Macro) \frac{\langle x, \Gamma \rangle \Downarrow v}{\langle (\mathbf{macro} \times), \Gamma \rangle \Downarrow \langle v, \Gamma(x) \in \mathbb{M} \rangle} \qquad \frac{\langle f, \Gamma \rangle \Downarrow \pi e \qquad \langle a, \Gamma \rangle \Downarrow i \qquad \langle i(e) \rangle \Downarrow \pi v}{\langle (\mathbf{f} \text{ a}), \Gamma \rangle \Downarrow \pi v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{m} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\langle f, \Gamma \rangle \Downarrow \pi e \quad \langle a, \Gamma \rangle \Downarrow i \quad \langle i(e) \rangle \Downarrow \pi v}{\langle (\mathbf{f} \text{ a}), \Gamma \rangle \Downarrow \pi v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{m} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow f \quad \langle f(e), \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad \langle m, \Gamma \rangle \Downarrow v}{\langle (\mathbf{f} \text{ e}), \Gamma \rangle \Downarrow v} \qquad (Apply-Impure) \\ \frac{\Gamma(m) \in \mathbb{M} \quad$$

Figure 2: The natural semantics of M.

2.1 Define

Def expressions are expressions of the form (def name value). They state that all symbols name evaluate to value, and evaluate to the the value they define. Multiple def expressions with the same name are invalid.

2.2 Symbol

Symbol expressions are expressions of the form name. They always evaluate to the value of the def expression that defines name. If name is not defined, they evaluate to \bot .

2.3 Function

Function expressions are expressions of the form (**fn** name value). They evaluate to a function λ name.value. When applied, they perform the substitution value [name/argument].

2.4 Impure

Impure expressions are expressions of the form (impure value) They evaluate to a process π value. When applied, they apply their argument to value. If the result of application is not a process, they evaluate to \bot instead.

2.5 Macro

Macro expressions are expressions of the form (macro name). They state that all symbols name are macros, and evaluate to the value of name. When applied, macros evaluate name and apply it to the expression encoding of argument, then evaluate the result.

2.6 Apply

Apply expressions are expressions of the form (value argument). If value is a function, it performs application as described in section 2.3. If value is a process, it performs application as described in section 2.4. If value is a macro, it performs application as described in section 2.5

3 Encodings

```
;;; Functions
                                               (def left?
                                                 (fn either
(def id (fn x x))
                                                   (ap either
                                                     (ap const true)
(def compose
                                                     (ap const false))))
  (fn f g
    (fn x
                                               (def right?
      (ap f g x))))
                                                 ({\bf fn}\ {\it either}\ 
                                                   (ap either
(def const
                                                     (ap const false)
  (fn x
                                                     (ap const true))))
    (fn "" x)))
                                               ;;; Lists
;;; Booleans
                                               (def nil (ap left false))
(def true
                                               (def cons (ap compose right pair))
  (fn x
    (fn "" x)))
                                               (def car
                                                 (fn value
(def false
                                                   (ap list id left)))
  (fn ""
    (fn x x)))
                                               (def cdr
                                                 (fn list
;;; Products
                                                   (ap list id right)))
(def pair
                                               (def nil? left?)
  (fn first second
    (fn value
                                               ;;; Natural Numbers
     (ap value first second))))
                                               (def 0 (ap left id))
                                               (def 0? left?)
(def first
                                               (def inc right)
  (fn pair
    (ap pair true)))
                                               (def dec
                                                 (fn nat
(def second
                                                   (ap nat left id)))
  (fn pair
   (ap pair false)))
                                               ;;; Chars
                                               (def nat->char id)
;;; Coproducts
                                               (def char->nat id)
(def left
  (fn value
                                              ;;; Symbols
    (fn first ""
                                               (def list->symbol id)
      (ap first value))))
                                               (def symbol->list id)
(def right
  (fn value
                                               ;;; Expressions
    (fn "" second
                                               (def symbol? left?)
      (ap second value))))
                                               (def list? right?)
```

4 Reference Implementation

4.1 Parser

()

4.2 Interpreter

()

4.3 Compiler

()