IF2130 – Organisasi dan Arsitektur Komputer

sumber: Greg Kesden, CMU 15-213, 2012

Representasi Informasi: Floating Point

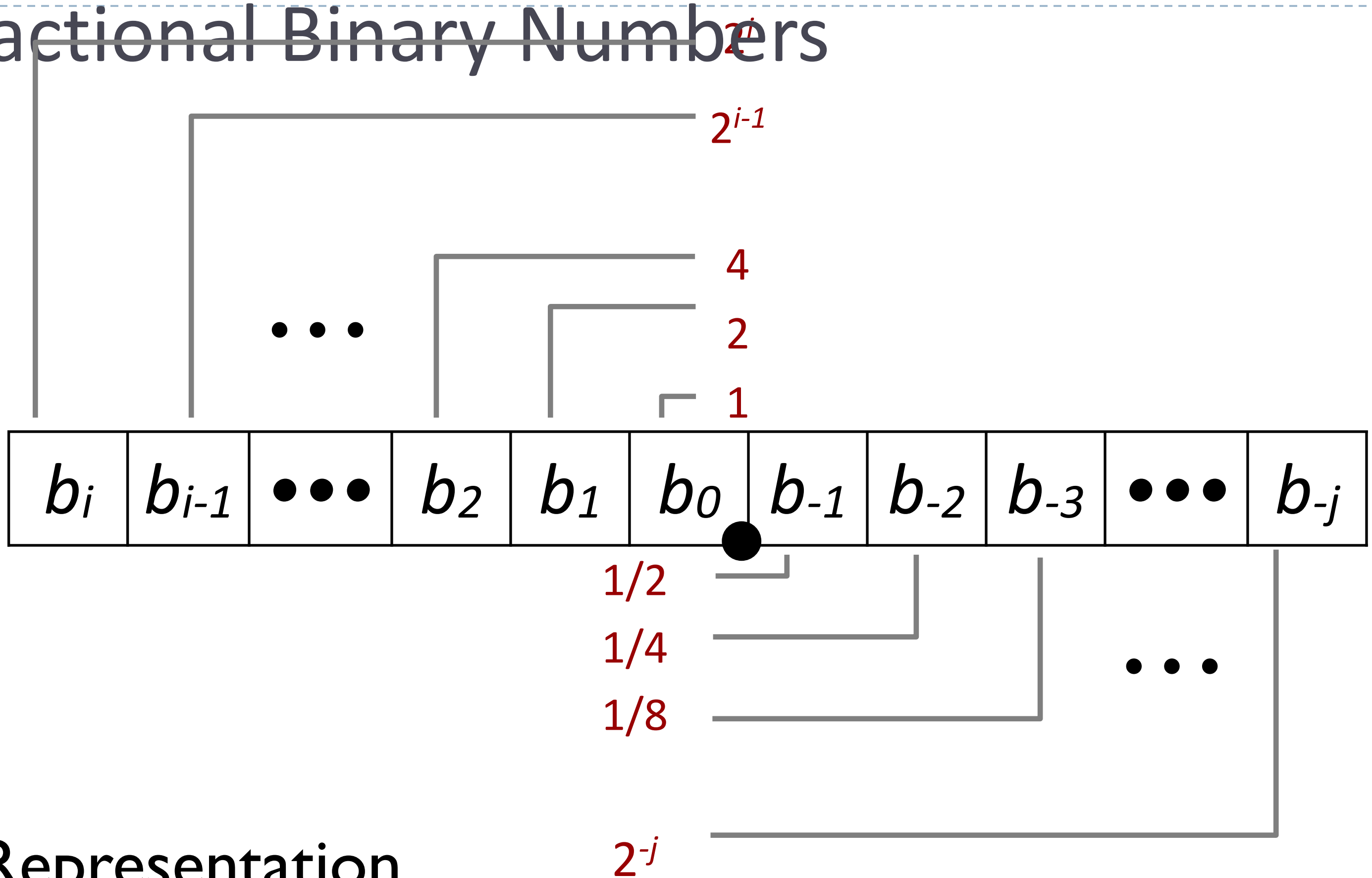
Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

Fractional binary numbers

- What is 1011.101_2?

Fractional Binary Numbers



- Representation
  - Bits to right of "binary point" represent fractional powers of 2
  - Represents rational number: $\sum_{k=-j}^{i} b_k \times 2^k$

Fractional Binary Numbers: Examples

- **Value** **Representation**

Value	Representation
5 3/4	101.11_2
2 7/8	10.111_2
1 7/16	1.0111_2

koma digeser ke kiri = dibagi 2

- **Observations**
  - Divide by 2 by shifting right (unsigned)
  - Multiply by 2 by shifting left
  - Numbers of form $0.111111\ldots_2$ are just below 1.0
    - $1/2 + 1/4 + 1/8 + \ldots + 1/2^i + \ldots \rightarrow 1.0$
    - Use notation $1.0 - \varepsilon$

i = jumlah bit

epsilon tergantung bit nya

Representable Numbers

- Limitation #1
  - Can only exactly represent numbers of the form $x/2^k$
    - Other rational numbers have repeating bit representations
  - Value — Representation
    - 1/3 — $0.0101010101[01]..._2$
    - 1/5 — $0.001100110011[0011]..._2$
    - 1/10 — $0.0001100110011[0011]..._2$

selain x/2^k harus dibulatkan/dipotong karena akan membentuk pattern dan berujung nilainya hanya mendekati, bukan nilai exact

- Limitation #2
  - Just one setting of decimal point within the *w* bits
    - Limited range of numbers (very small values? very large?)

kalau nilai terlalu besar, gabisa membentuk pecahan yang kecil

Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

IEEE Floating Point

- IEEE Standard 754
  - Established in 1985 as uniform standard for floating point arithmetic
    - Before that, many idiosyncratic formats
  - Supported by all major CPUs

- Driven by numerical concerns
  - Nice standards for rounding, overflow, underflow
  - Hard to make fast in hardware
    - Numerical analysts predominated over hardware designers in defining standard

Floating Point Representation

- Numerical Form:

M = Mantissa

$$(-1)^s \, M \; 2^E$$

sign bit 1 = (-)
sign bit 0 = (+)

  - **Sign bit *s*** determines whether number is negative or positive
  - **Significand *M*** normally a fractional value in range [1.0,2.0).
  - **Exponent *E*** weights value by power of two

- Encoding
  - MSB s is sign bit ***s***
  - exp field encodes ***E*** (but is not equal to E)
  - frac field encodes ***M*** (but is not equal to M)

s	exp	frac

E dikodekan jadi bit di exp

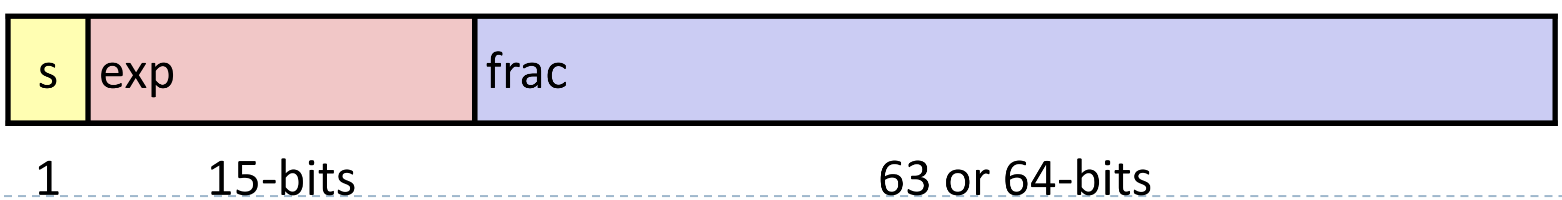
M dikodekan jadi bit di frac

Precision options

- Single precision: 32 bits

- Double precision: 64 bits

- Extended precision: 80 bits (Intel only)



"Normalized" Values

- When: exp ≠ 000…0 and exp ≠ 111…1

rentang M adalah 1 sampai 2 (1 <= M < 2)

- Exponent coded as a ***biased*** value: ***E = Exp – Bias***
  - *Exp*: unsigned value exp
  - *Bias* = 2^{k-1} - 1, where *k* is number of exponent bits
    - Single precision: 127 (Exp: 1…254, E: -126…127)
    - Double precision: 1023 (Exp: 1…2046, E: -1022…1023)

E rentang dr - sampai + agar bisa membentuk pecahan dan bilangan yg sgt besar

- Significand coded with implied leading 1: ***M*** = $1.xxx…x_2$
  - xxx…x: bits of frac
  - Minimum when frac=000…0 (*M* = 1.0)
  - Maximum when frac=111…1 (*M* = 2.0 – ε)
  - Get extra leading bit for "free"

Normalized Encoding Example

- Value: `Float F = 15213.0;`
  - $15213_{10} = 11101101101101_2$

$$= 1.\underbrace{1101101101101_2}_{\text{frac}} \times 2^{\overset{E}{13}}$$

geser titiknya 13 ke kiri

- Significand

$M = 1.\underline{1101101101101}_2$

frac= $\underline{1101101101101}0000000000_2$

- Exponent

$E = 13$

$Bias = 127$

$Exp = 140 = 10001100_2$

exp = E + bias = 13 + 127 = 140

nilai terkecil = 1 - bias = 1 - 127 = -126

- Result:

32 bit (single precision)

s	exp	frac
0	10001100	11011011011010000000000

Denormalized Values

- Condition: exp = 000…0

leading 0

- Exponent value: ***E* = –*Bias* + 1** (instead of ***E*** = 0 – ***Bias***)
- Significand coded with implied leading 0: ***M*** = $0.xxx\ldots x_2$
  - `xxx…x`: bits of `frac`
- Cases

0, sehingga rentangnya tuh makin kecil

  - `exp` = 000…0, `frac` = 000…0
    - Represents zero value
    - Note distinct values: +0 and –0 (why?)

ketika exp = 0 dan frac = 0, maka nilainya 0

  - `exp` = 000…0, `frac` ≠ 000…0
    - Numbers closest to 0.0
    - Equispaced

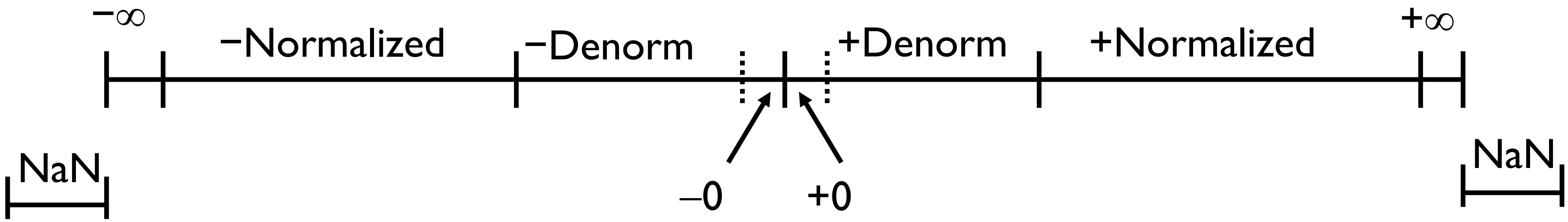
Special Values

- Condition: **`exp`** = 111…1

- Case: **`exp`** = 111…1, **`frac`** = 000…0
  - Represents value ∞ (infinity)
  - Operation that overflows
  - Both positive and negative
  - E.g., $1.0/0.0 = -1.0/-0.0 = +\infty$, $1.0/-0.0 = -\infty$

- Case: **`exp`** = 111…1, **`frac`** ≠ 000…0
  - Not-a-Number (NaN)
  - Represents case when no numeric value can be determined
  - E.g., sqrt(–1), $\infty - \infty$, $\infty \times 0$

Visualization: Floating Point Encodings



Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

Tiny Floating Point Example

s	exp	frac
1	4-bits	3-bits

bias = 2^k-1 - 1
= 2^4-1 -1
= 2^3 - 1
= 7

- 8-bit Floating Point Representation
  - the sign bit is in the most significant bit
  - the next four bits are the exponent, with a bias of 7
  - the last three bits are the `frac`

rentang dari -6 sampai 7

denormalize
E = -bias +1 = -7+1 = -6

- Same general form as IEEE Format
  - normalized, denormalized
  - representation of 0, NaN, infinity

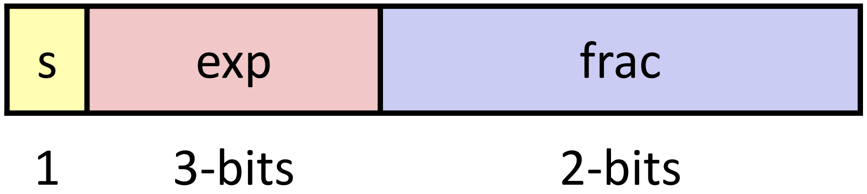
normalize
Exp = E - bias
4 = E - 7

Dynamic Range (Positive Only)

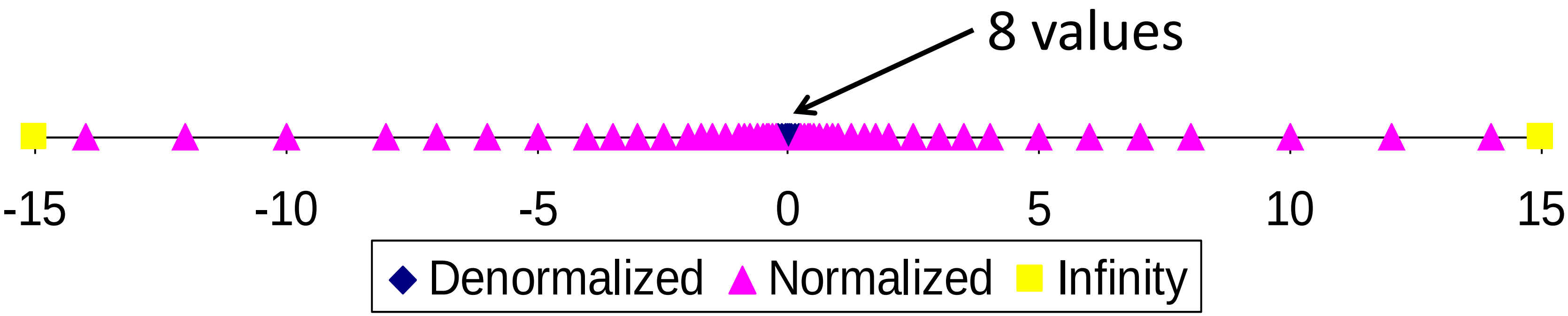
	s	exp	frac	E	Value	
			2^n		2^n x 2^E	
	0	0000	000	-6	0	
	0	0000	001	-6	1/8*1/64 = 1/512	**closest to zero**
Denormalized numbers	0	0000	010	-6	2/8*1/64 = 2/512	
	…					denor E = -bias+1
	0	0000	110	-6	6/8*1/64 = 6/512	
	0	0000	111	-6	7/8*1/64 = 7/512	**largest denorm**
	0	0001	000	-6	8/8*1/64 = 8/512	**smallest norm**
normalized	0	0001	001 1 - 7 = -6	-6	9/8*1/64 = 9/512	
	…	exp				normalize E = Exp - bias
	0	0110	110 6 - 7 = -1	-1	14/8*1/2 = 14/16	
	0	0110	111	-1	15/8*1/2 = 15/16	**closest to 1 below**
Normalized numbers	0	0111	000 7 - 7 = 0	0	8/8*1 = 1	
	0	0111	001	0	9/8*1 = 9/8	**closest to 1 above**
	0	0111	010	0	10/8*1 = 10/8	
	…		14 - 7 = 7			
	0	1110	110	7	14/8*128 = 224	
	0	1110	111	7	15/8*128 = 240	**largest norm**
	0	1111	000	n/a	inf	
		standar dr infinity				

Distribution of Values

- 6-bit IEEE-like format
  - e = 3 exponent bits
  - f = 2 fraction bits
  - Bias is $2^{3-1}-1 = 3$

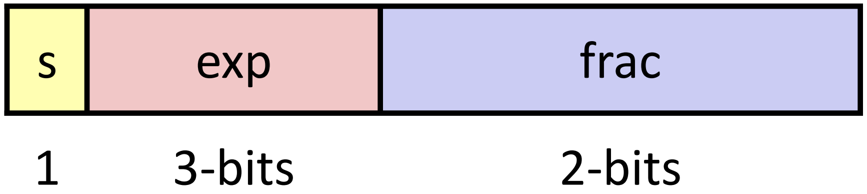


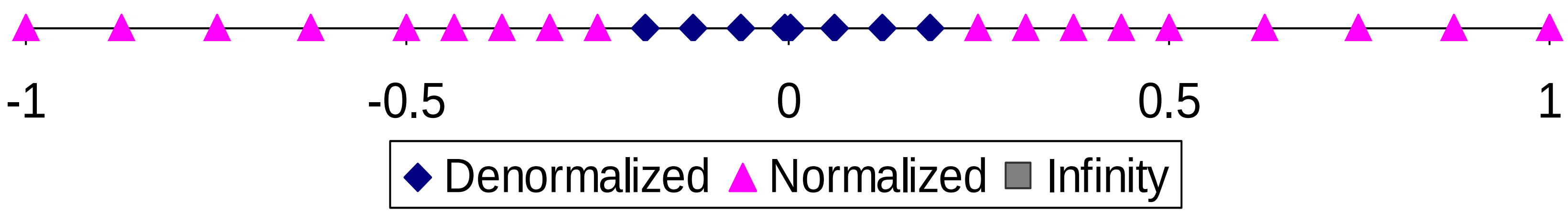
- Notice how the distribution gets denser toward zero.



Distribution of Values (close-up view)

- 6-bit IEEE-like format
  - e = 3 exponent bits
  - f = 2 fraction bits
  - Bias is 3





Special Properties of the IEEE Encoding

- FP Zero Same as Integer Zero
  - All bits = 0
- Can (Almost) Use Unsigned Integer Comparison
  - Must first compare sign bits
  - Must consider $-0 = 0$
  - NaNs problematic
    - Will be greater than any other values
    - What should comparison yield?
  - Otherwise OK
    - Denorm vs. normalized
    - Normalized vs. infinity

Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

Floating Point Operations: Basic Idea

- $x +_f y = \mathrm{Round}(x + y)$
- $x \times_f y = \mathrm{Round}(x \times y)$
- Basic idea
  - First compute exact result
  - Make it fit into desired precision
    - Possibly overflow if exponent too large
    - Possibly round to fit into `frac`

Rounding

- Rounding Modes (illustrate with $ rounding)

	$1.40	$1.60	$1.50	$2.50	–$1.50
Towards zero	$1	$1	$1	$2	–$1
Round down (−∞)	$1	$1	$1	$2	–$2
Round up (+∞)	$2	$2	$2	$3	–$1
Nearest Even(default)	$1	$2	$2	$2	–$2

Closer Look at Round-To-Even

- Default Rounding Mode
  - Hard to get any other kind without dropping into assembly
  - All others are statistically biased
    - Sum of set of positive numbers will consistently be over- or under-estimated
- Applying to Other Decimal Places / Bit Positions
  - When exactly halfway between two possible values
    - Round so that least significant digit is even
  - E.g., round to nearest hundredth

1.2349999	1.23	(Less than half way)
1.2350001	1.24	(Greater than half way)
1.2350000	1.24	(Half way—round up)
1.2450000	1.24	(Half way—round down)

nearest even
=
ke genap yang paling deket

Rounding Binary Numbers

- Binary Fractional Numbers
  - "Even" when least significant bit is 0
  - "Half way" when bits to right of rounding position = $100\ldots_2$

kalau misal pas/lebih dr 1/2 -> kalau depannya 0 jd dibulatin ke bawah
-> kalau depannya 1 jd dibulatin ke atas

kalau misal kurang dr 1/2 -> mau dpnnya 0/1 ttp bulatin ke bawah

- Examples
  - Round to nearest 1/4 (2 bits right of binary point)

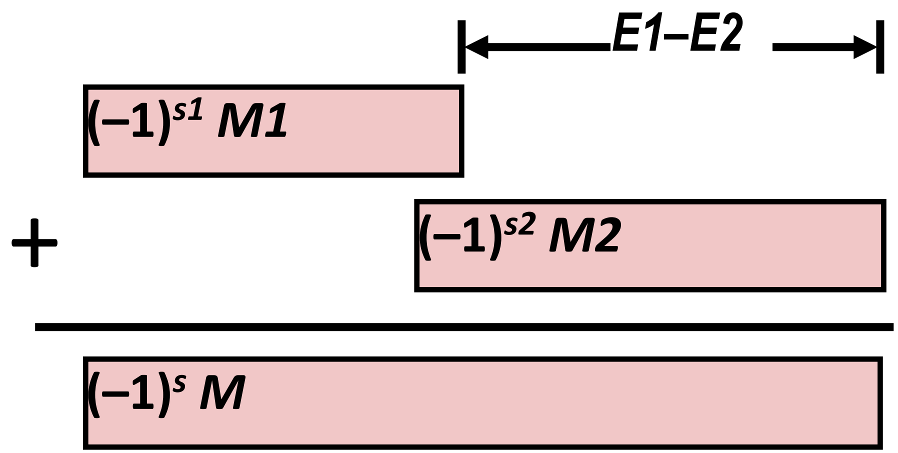
Value	Binary	Rounded	Action	Rounded Value
2 3/32	10.00011_2	10.00_2	(<1/2—down)	2
2 3/16	10.00110_2	10.01_2	(>1/2—up)	2 1/4
2 7/8	10.11100_2	11.00_2	(1/2—up)	3
2 5/8	10.10100_2	10.10_2	(1/2—down)	2 1/2

FP Multiplication

- $(-1)^{s1}\ M1\ \ 2^{E1}\ \ \times\ \ (-1)^{s2}\ M2\ \ 2^{E2}$
- Exact Result: $(-1)^{s}\ M\ \ 2^{E}$
  - Sign s: $s1$ ^ $s2$
  - Significand M: $M1 \times M2$
  - Exponent E: $E1 + E2$
- Fixing
  - If $M \geq 2$, shift M right, increment E
  - If E out of range, overflow
  - Round M to fit `frac` precision
- Implementation
  - Biggest chore is multiplying significands

Floating Point Addition

- $(-1)^{s1}\ M1\ \ 2^{E1}\ \ +\ \ (-1)^{s2}\ M2\ \ 2^{E2}$
  - Assume $E1 > E2$
- Exact Result: $(-1)^{s}\ M\ \ 2^{E}$
  - Sign s, significand M:
    - Result of signed align & add
  - Exponent E: $E1$



- Fixing
  - If $M \geq 2$, shift M right, increment E
  - if $M < 1$, shift M left k positions, decrement E by k
  - Overflow if E out of range
  - Round M to fit `frac` precision

Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

Floating Point in C

- C Guarantees Two Levels
  - **`float`** single precision
  - **`double`** double precision
- Conversions/Casting
  - Casting between **`int`**, **`float`**, and **`double`** changes bit representation
  - **`double`/`float`** → **`int`**
    - Truncates fractional part
    - Like rounding toward zero
    - Not defined when out of range or NaN: Generally sets to TMin
  - **`int`** → **`double`**
    - Exact conversion, as long as **`int`** has ≤ 53 bit word size
  - **`int`** → **`float`**
    - Will round according to rounding mode

Summary

- IEEE Floating Point has clear mathematical properties
- Represents numbers of form M x 2^E
- One can reason about operations independent of implementation
  - As if computed with perfect precision and then rounded
- Not the same as real arithmetic
  - Violates associativity/distributivity
  - Makes life difficult for compilers & serious numerical applications programmers

Floating Point Puzzles

- For each of the following C expressions, either:
  - Argue that it is true for all argument values
  - Explain why not true

```
int x = ...;
float f = ...;
double d = ...;
```

Assume neither **d** nor **f** is NaN

- x == (int)(float) x
- x == (int)(double) x
- f == (float)(double) f
- d == (float) d
- f == -(-f);
- 2/3 == 2/3.0
- d < 0.0 ⇒ ((d*2) < 0.0)
- d > f ⇒ -f > -d
- d * d >= 0.0
- (d+f)-d == f

Interesting Numbers `{single,double}`

Description	*exp*	*frac*	*Numeric Value*
▸ Zero	00…00	00…00	0.0
▸ Smallest Pos. Denorm.	00…00	00…01	$2^{-\{23,52\}} \times 2^{-\{126,1022\}}$
▸ Single ≈ 1.4×10^{-45}			
▸ Double ≈ 4.9×10^{-324}			
▸ Largest Denormalized	00…00	11…11	$(1.0 - \varepsilon) \times 2^{-\{126,1022\}}$
▸ Single ≈ 1.18×10^{-38}			
▸ Double ≈ 2.2×10^{-308}			
▸ Smallest Pos. Normalized	00…01	00…00	$1.0 \times 2^{-\{126,1022\}}$
▸ Just larger than largest denormalized			
▸ One	01…11	00…00	1.0
▸ Largest Normalized	11…10	11…11	$(2.0 - \varepsilon) \times 2^{\{127,1023\}}$
▸ Single ≈ 3.4×10^{38}			
▸ Double ≈ 1.8×10^{308}			

Creating Floating Point Number

- Steps
  - Normalize to have leading 1
  - Round to fit within fraction
  - Postnormalize to deal with effects of rounding

s	exp	frac
1	4-bits	3-bits

- Case Study
  - Convert 8-bit unsigned numbers to tiny floating point format

Example Numbers

128	10000000
15	00001111
17	00010001
19	00010011
33	00100001
35	00100011
138	10001010
63	00111111

Rounding

1.BBGRXXX

Guard bit: LSB of result

Sticky bit: OR of remaining bits

Round bit: 1st bit removed

- Round up conditions
  - Round = 1, Sticky = 1 → > 0.5
  - Guard = 1, Round = 1, Sticky = 0 → Round to even

Value	***Fraction***	***GRS***	***Incr?***	***Rounded***
128	1.0000000	000	N	1.000
15	1.1110000	100	N	1.111
17	1.0001000	010	N	1.000
19	1.0011000	110	Y	1.010
138	1.0001010	011	Y	1.001
63	1.1111100	111	Y	10.000

More Slides

Today: Floating Point

- Background: Fractional binary numbers
- IEEE floating point standard: Definition
- Example and properties
- Rounding, addition, multiplication
- Floating point in C
- Summary

Mathematical Properties of FP Add

- Compare to those of Abelian Group
  - Closed under addition?
    - But may generate infinity or NaN
  - Commutative?
  - Associative?
    - Overflow and inexactness of rounding
  - 0 is additive identity?
  - Every element has additive inverse
    - Except for infinities & NaNs
- Monotonicity
  - $a \geq b \Rightarrow a+c \geq b+c$?
    - Except for infinities & NaNs

Mathematical Properties of FP Mult

- Compare to Commutative Ring
  - Closed under multiplication?
    - But may generate infinity or NaN
  - Multiplication Commutative?
  - Multiplication is Associative?
    - Possibility of overflow, inexactness of rounding
  - 1 is multiplicative identity?
  - Multiplication distributes over addition?
    - Possibility of overflow, inexactness of rounding

- Monotonicity
  - $a \geq b \;\; \& \; c \geq 0 \;\; \Rightarrow a * c \geq b * c$?
    - Except for infinities & NaNs

Normalize

s	exp	frac
1	4-bits	3-bits

- Requirement
  - Set binary point so that numbers of form 1.xxxxx
  - Adjust all to have leading one
    - Decrement exponent as shift left

Value	*Binary*	*Fraction*	*Exponent*
128	10000000	1.0000000	7
15	00001101	1.1010000	3
17	00010001	1.0001000	4
19	00010011	1.0011000	4
138	10001010	1.0001010	7
63	00111111	1.1111100	5

Postnormalize

- Issue
  - Rounding may have caused overflow
  - Handle by shifting right once & incrementing exponent

Value	*Rounded*	*Exp*	*Adjusted*	*Result*
128	1.000	7		128
15	1.101	3		15
17	1.000	4		16
19	1.010	4		20
138	1.001	7		134
63	10.000	5	1.000/6	64