

# NFA dgn $\epsilon$ -Transition

### **NFA's With €-Transitions**



- We can allow state-to-state transitions on ∈ input.
- These transitions are done spontaneously, without looking at the input string.
- A convenience at times, but still only regular languages are accepted.

#### **Contoh Soal**



#### Soal

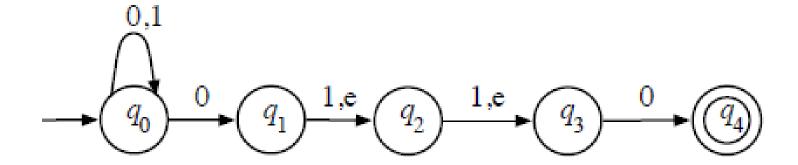
An NFA for the language of all strings over  $\{0,1\}$  that end with one of  $0110,\,010,\,$  and 00.

An NFA for the language of all strings over  $\{a, b, c\}$  for which one of (the number of occurrences of a), (the number of occurrences of b), and (the number of occurrences of c) is a multiple of a.



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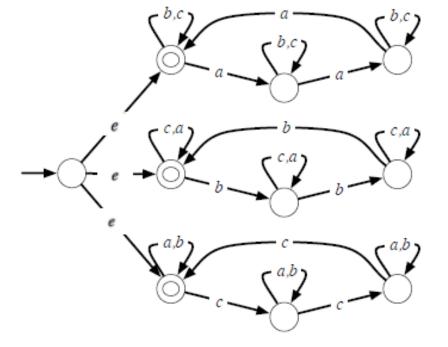
#### Jawab





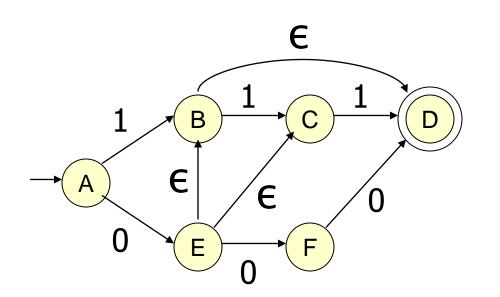
An NFA for the language of all strings over  $\{a, b, c\}$  for which one of (the number of occurrences of a), (the number of occurrences of b), and (the number of occurrences of c) is a multiple of a.

Jawab



# **NFA** dgn $\epsilon$ -Transition





$$\begin{array}{c|cccc} & 0 & 1 & \boldsymbol{\epsilon} \\ \hline \rightarrow & A & \{E\} & \{B\} & \varnothing \\ & B & \varnothing & \{C\} & \{D\} \\ & C & \varnothing & \{D\} & \varnothing \\ & * & D & \varnothing & \varnothing \\ & * & D & \varnothing & \varnothing \\ & E & \{F\} & \varnothing & \{B, C\} \\ & F & \{D\} & \varnothing & \varnothing \end{array}$$

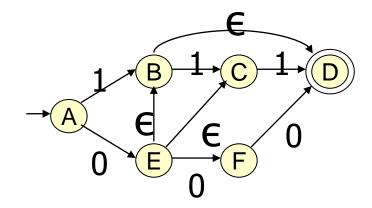
 memungkinkan adanya transisi antar state dgn input ε

#### **Closure of States**



 CL(q) = sekumpulan state yang bisa dicapai dari state q dengan hanya mengikuti busur berlabel ε termasuk status dirinya.

- CL(E) = {B, C, D, E}
- Closure dari kumpulan state:
  - gabungan closure setiap state



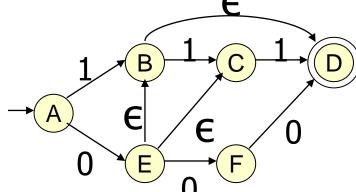
#### **Extended Delta**



- Basis:  $\delta(q, \epsilon) = CL(q)$ .
- Induction:  $\delta(q, xa)$  sbb:
  - 1. misal  $\delta(q, x) = S$ .
  - 2. = union dari  $CL(\delta(p, a))$  utk semua p di S.
- Intuisi:  $\delta(q, w)$  adalah sekumpulan state yang dapat dicapai dari q dengan mengikuti path busur w



## **Cth: Extended Delta**



- $\delta(A, \epsilon) = CL(A) = \{A\}.$
- $\delta(A, 0) = CL(\{E\}) = \{B, C, D, E\}.$
- $\delta(A, 01) = CL(\{C, D\}) = \{C, D\}.$
- Language dari sebuah  $\epsilon$ -NFA adalah sekumpulan string w dimana  $\delta(q_0, w)$  mengandung final state.