

Algorithms: Design and Analysis, Part II

# Huffman Codes

Introduction and Motivation

## Binary Codes

Binary code: Maps each character of an alphabet  $\Sigma$  to a binary string.

Example:  $\Sigma = \text{a-z}$  and various punctuation (size 32 overall, say)

Obvious encoding: Use the 32 5-bit binary strings to encode this  $\Sigma$  (a fixed-length code)

Can we do better? Yes, if some characters of  $\Sigma$  are much more frequent than others, using a variable-length code.

## **Ambiguity**

Example: Suppose  $\Sigma = \{A,B,C,D\}$ . Fixed-length encoding would be  $\{00,01,10,11\}$ .

Suppose instead we use the encoding  $\{0,01,10,1\}$ . What is 001 an encoding of?

- A) AB  $\rightarrow$  Leads to 001
- B) CD
- C) AAD  $\rightarrow$  Also leads to 001
- D) Not enough info to answer question

### Prefix-Free Codes

Problem: With variable-length codes, not clear where one character ends + the next one begins.

Solution: Prefix-free codes - make sure that for every pair  $i, j \in \Sigma$ , neither of the encodings f(i), f(j) is a prefix of the other.

Example: {0,10,110,111}

Why useful? Can give shorter encodings with non-uniform character frequencies.

## Example

#### Example:

Α	60%	00	0
В	25%	01	10
C	10%	10	110
D	5%	11	111

 $\Sigma$  frequencies fixed-length variable-length (prefix free)

Fixed-length encoding: 2 bits/character

Variable-length encoding: How many bits needed on average?

$$0.6 \cdot 1 + 0.25 \cdot 2 + (0.1 + 0.05) \cdot 3 = 1.55$$