bpfbox: Simple Precise Process Confinement with eBPF

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Outline of Talk

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- 2. The Status Quo
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- 5. bpfbox Design and Implementation
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The Process Confinement Problem

What is Process Confinement?

We want to be able to **confine** our **processes**.

► Also known as sandboxing

Why do we want to do this?

- ▶ Default protection mechanisms are too:
 - ► Granular
 - ► User-centric
 - Discretionary
- ► Protection can be **overridden**
 - ► Superuser (*nix)
 - ► Administrator (Windows)

How do we protect the user from themselves and their own processes?

An Interesting Question

How many processes do you think are running on your computer **right now**?

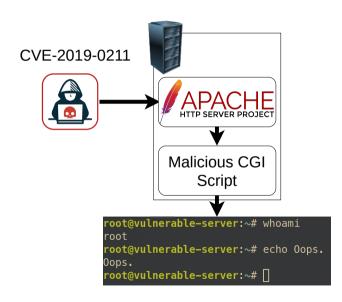
- ► Probably a lot more than you think
- ► You probably didn't start most of them yourself
- ► You might not even know what some of them are for

Compromised Processes

- Web servers
- ▶ Daemons
- Chat applications
- etc.

The Morris Worm

- ► Backdoor in sendmail daemon
- ► Buffer overflow in fingerd
- ► Estimated damage: \$100,000-\$10,000,000



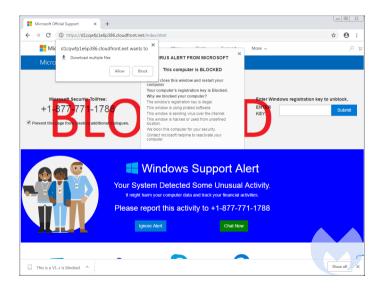
Semi-Honest Software

- ► Software that does its job...
- ► But also performs potentially unwanted actions



Malicious Software

- Viruses
- ▶ Trojans
- ► Ransomware
- ► Spyware
- etc.



Attack goals?

- ► Installing backdoors/rootkits
- ► Information leakage
- ▶ Denial of service
- ▶ Data ransom
- ► Setting up a botnet

Process confinement reduces the attack surface.

The Process Confinement Problem

► "A Note on the Confinement Problem" (Lampson, 1973)

Systems

C. Weissman
Editor

A Note on the
Confinement Problem

Butler W. Lampson Xerox Palo Alto Research Center

► An open problem for nearly **five decades**

The Status Quo

Unix DAC

POSIX Capabilities

Namespaces and Cgroups

System Call Interposition

Linux MAC

Containers / Containerized Package Management

eBPF 101

eBPF in the Beginning

eBPF Extended Berkley Packet Filter

- ▶ But it has little to do with Berkley, packets, or filtering nowadays
- ► The name BPF is preserved for historical reasons

So then what is eBPF?

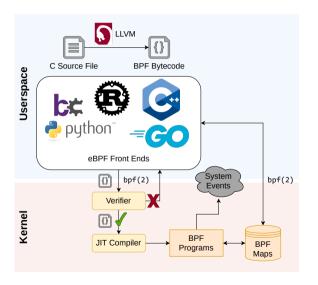
- ► A major re-write of the Linux BPF engine
 - ► Alexei Starovoitov and Daniel Borkmann
- ► Merged into the Linux kernel in 2014
- ► The original goal was fine-grained, cross-layer system introspection

What Can eBPF Do?



Direct Packet Userspace Hardware LSM Hooks **Functions** Access **Block Device System Calls** TCP / IP I/O And much more... Kernelspace Sockets Stack Traces **Functions**

How eBPF Works



Verifiably Safe Programs

► TODO explain verifier

eBPF in 2020

eBPF is now more than just an observability tool.

- ► eBPF provides a **safe**, **efficient**, and **flexible** way for privileged users to extend the kernel
- ► eBPF turns Linux into a **programmable kernel**

Linux $5.7 \rightarrow KRSI$ (Kernel Runtime Security Instrumentation)

- ► Attach BPF programs to LSM hooks
- ► Make security decisions and generate audit logs with eBPF

KRSI: BPF LSM Programs

► TODO explain KRSI with a picture

bpfbox Overview

bpfbox at a Glance

- bpfbox is a novel process confinement mechanism for Linux
 - ► Using a new Linux technology called eBPF

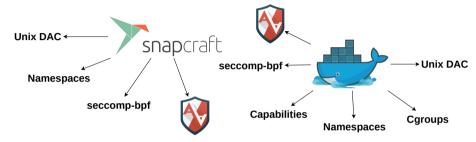
 Users write per-application policy in a simple policy language

- ▶ Policy is enforced by attaching eBPF programs to LSM hooks
 - ► Integrates cross-layer state into policy decisions



Motivation

► Existing process confinement mechanisms are **complex**



Existing process confinement mechanisms are difficult to use



► Can we do any better?

eBPF Changes the Game

eBPF enables:

- ► Fine-grained system introspection
- ► Integration of **cross-layer state** with policy enforcement
- ► Rapid prototyping
- ► Safe production deployment of new security solutions

We have an opportunity to rethink process confinement from the ground up.

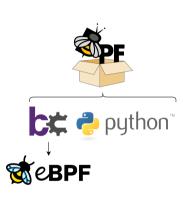
bpfbox Design and Implementation

bpfbox Implementation

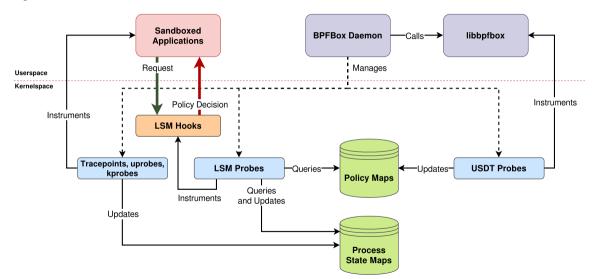
► Userspace daemon using the Python3 bcc framework

- ► Kernelspace components are all written in eBPF
 - ► LSM probes (KRSI), kprobes, uprobes
 - ► Under 2000 source lines of kernelspace code

- ► Thanks to eBPF, bpfbox is **light-weight**, **flexible**, and **production-safe**
 - ► Works out of the box on any vanilla Linux kernel > 5.8



bpfbox Architecture



Policy Design Goals

1. Simplicity

► Policy should be simple enough for ad hoc confinement

2. Application transparency

▶ Policy should not require changes to the confined application

3. Flexibility

► Policy should offer optional layers of granularity

4. Security

- ► Policy should follow the principle of least privilege
- ► It should be difficult to write an insecure policy

Rules and Directives

Rules specify access to system objects:

- ► fs(file, access)
- ▶ net(socket, access)
- ► signal(prog, sig)
- etc.

Directives augment blocks of rules:

- ► #[directive] syntax
- ► Specify actions to be taken on a block of rules
- ► Add additional context to a block of rules

Taints and Transitions

- ▶ #[taint] → Start confinement
- ▶ #[transition] → Switch profiles on execve

```
#![profile "/bin/mywebdaemon"]
#[taint] {
    net(inet, any)
    net(inet6, any)
}
/* ... */
#[transition] {
    fs("/bin/myhelper", getattr|read|exec)
}
```

Policy at the Function Call Level

- ▶ #[func "foo"] → Apply rules only within a call to foo()
- ▶ #[kfunc "foo"] → Same thing, but for kernel functions

```
#![profile "/sbin/mylogin"]
#[func "check_password"]
#[allow] {
    fs("/etc/passwd", read)
    fs("/etc/shadow", read)
#[func "add_user"]
#[allow] {
    fs("/etc/passwd", read|append)
    fs("/etc/shadow", read|append)
/* ... */
```

bpfbox Performance Evaluation

Methodology

- ► Phoronix Test Suite OSBench
 - ► Measures basic OS functionality
 - ► (spawning processes, memory allocations, etc.)
- ► Phoronix Test Suite Apache
 - ► Benchmark Apache httpd packets per second
- ► Kernel compilation benchmarks
 - ► Measure Linux kernel compilation performance
 - ► Heavy workload, spawning lots of processes

Methodology

Two modes of operation for each test.

- ► Passive mode
 - bpfbox and AppArmor instrument hooks, but do not enforce or audit
 - ► Test lowest possible overhead

- ► Complaining mode
 - ▶ bpfbox and AppArmor complain about (log) every security-sensitive operation
 - ► Test worst case overhead

Results

- ► Phoronix OSBench
 - ► Passive: bpfbox is **roughly equivalent** to AppArmor
 - ► Complaining: bpfbox performs **significantly better** than AppArmor
- ▶ Phoronix Apache
 - bpfbox and AppArmor are roughly equivalent
- ► Kernel compilation
 - ► Passive: bpfbox is **roughly equivalent** to AppArmor
 - Complaining: bpfbox performs better in kernelspace overhead and worse in userspace overhead

The Future of eBPF and Security

Conclusion

Acknowledgements

Special thanks to:

- ► Alexei Starovoitov and Daniel Borkmann (creators of eBPF)
- ► K.P. Singh (creator of KRSI)
- Fellow **bcc contributors** (an awesome eBPF framework)
- ► Anonymous CCSW'2020 reviewers (valuable feedback)

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Contributions

- ► First policy enforcement engine written in eBPF
- ► Integration of userspace and kernelspace state with LSM layer enforcement
- ► A simple policy language for ad hoc process confinement
 - ► But with optional complexity for **fine-grained protection**



github.com/willfindlay/bpfbox Check out the project on GitHub!