



Host-Based Anomaly Detection with Extended BPF

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eBPF at a High Level

- ◆ Inject user-specified code into the kernel
- ◆ BPF code runs in kernelspace, can instrument essentially **all system behavior**
- ◆ This sounds a lot like a kernel module...
 - Key difference? **Safety.**
- ◆ Before they can run in the kernel, BPF programs are **statically verified**

◆ **Performance monitoring**

- Netflix
- Facebook
- Google
- ... many others

◆ **Established tools**

- bcc-tools (over 100 performance monitoring / visibility tools)

◆ **Network security**

- Cloudflare's DDoS mitigation stack

- ◆ **A lot of security is about what we can see**
 - eBPF lets you see **everything** about your system
 - ... and it can do this with **crazy low overhead**
- ◆ **Before eBPF, system introspection came at a cost**
 - Speed
 - Scope
 - Production safety
- ◆ **eBPF can do everything, without the speed / scope / safety trade-off**
 - Although eBPF comes with its own nuances (more on this later)

eBPF Architecture

Canada's Capital University

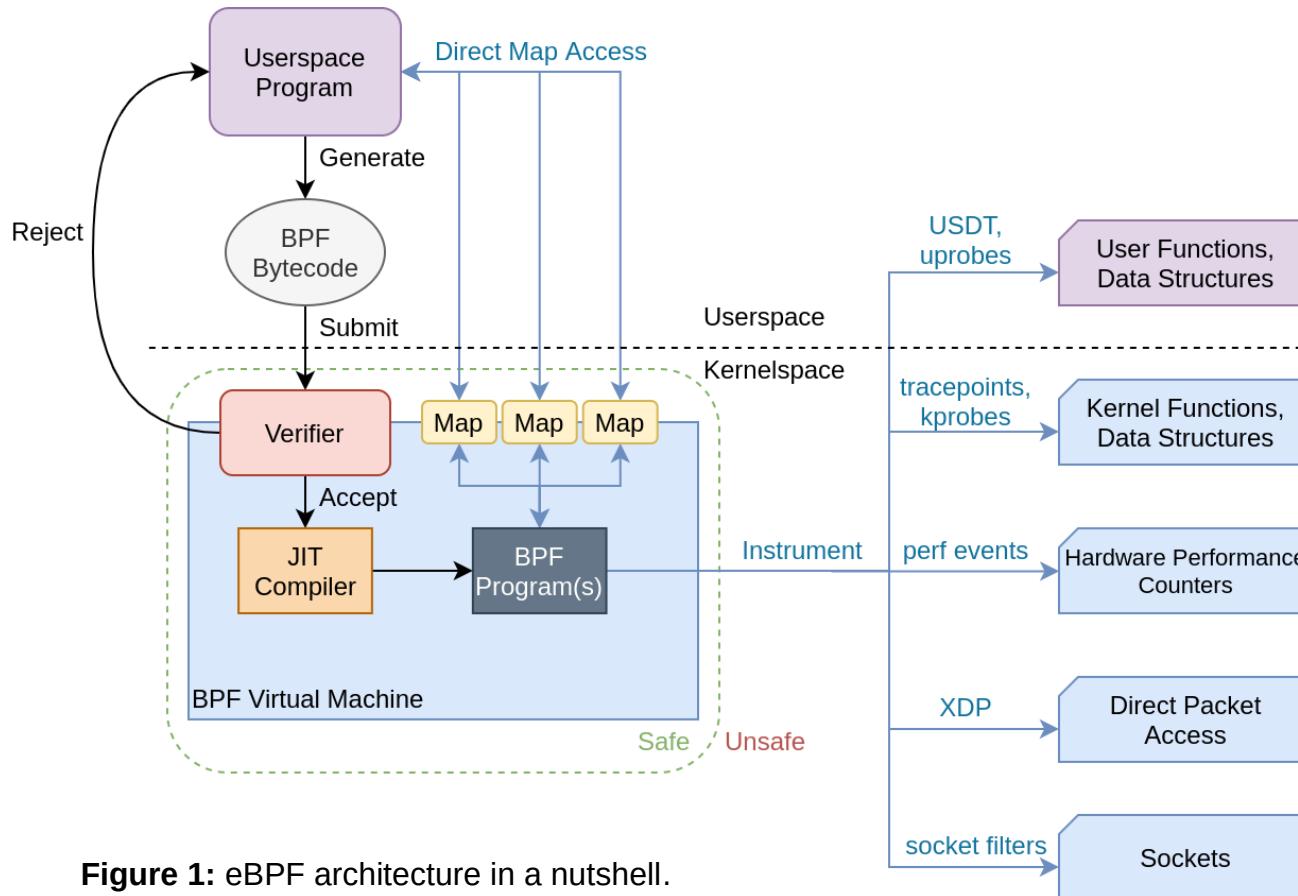


Figure 1: eBPF architecture in a nutshell.
Note that the list of program types is not exhaustive.

◆ eBPF verifier

- Ensure BPF program will **not** crash the kernel
- 10,000 lines of C code in kernel
- BPF system call traps to verifier on PROG_LOAD

◆ How to guarantee safety?

Limitations + simulation + static analysis

- **512 byte** stack space
- No **unbounded loops**
- Max **1 million BPF instructions** per program
- No buffer access with **unbounded induction variable**
- Etc.

BPF Programs Still Can Be Complex

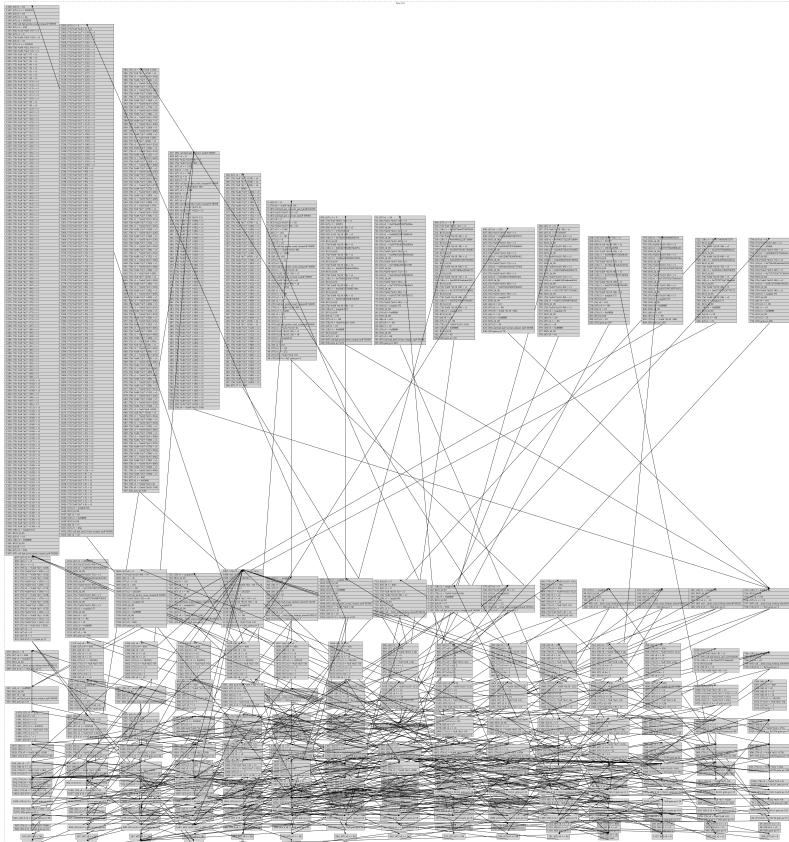


Figure 2: Instruction flow graph of ebpH's sys_exit tracepoint.

- ◆ **ebpH's sys_exit tracepoint**
 - bpftrace + graphviz usage
 - 1,574 BPF instructions
 - 1,930 machine instructions
- ◆ **BPF programs can interact with each other**
 - Direct map access
 - Tail calls

- ◆ **Early anomaly detection system by Anil Somayaji**
- ◆ **The idea:**
 - Instrument system calls to build per-executable behavioral profiles
 - Delay anomalous system calls proportionally to recent anomalies
- ◆ **Problems?**
 - Implemented as a **kernel patch**
 - Need to make crazy modifications for it to work
 - Patch the scheduler, write in assembly language, etc.
 - Not production-safe
 - Not portable

ebpH: Back to the Future

◆ ebpH

- “Extended BPF + Process Homeostasis”
- 20 year old technology...
- Re-written using modern technology

Table 1: Comparing ebpH and pH.

System	Implementation	Portable	Production Safe	Low Mem. Overhead	Low Perf. Overhead	Detection	Response
pH	Kernel Patch	✗	✗	✓	✓	✓	✓
ebpH	eBPF + Userspace Daemon	✓	✓	✗	✓	✓	✗

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◆ ebpH

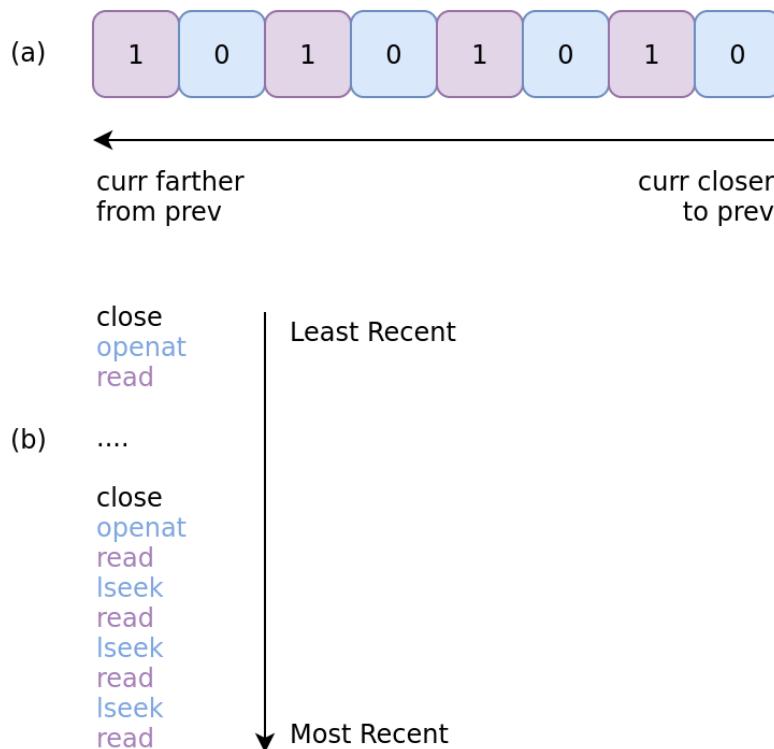
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ebpH in Detail

```
binary = ls  
(curr, prev) = (read, close)
```



Same idea as pH:

- ◆ Trace system calls
 - ◆ Build profile of lookahead pairs
 - ◆ Gather enough data
 - ◆ Flag new lookahead pairs as anomalies

eBPF makes this **safe**.

How ebpH Collects Data

- ◆ **Tracepoints (static kernel tracing)**
 - Instrument **system calls**
 - Instrument **scheduler**
- ◆ **Kprobes (dynamic kernel tracing)**
 - Instrument **signal delivery**
- ◆ **Uprobes (dynamic user tracing)**
 - Instrument **libebph.so**
 - Allow user to issue commands to ebpH's BPF programs

ebpH Architecture

Canada's Capital University

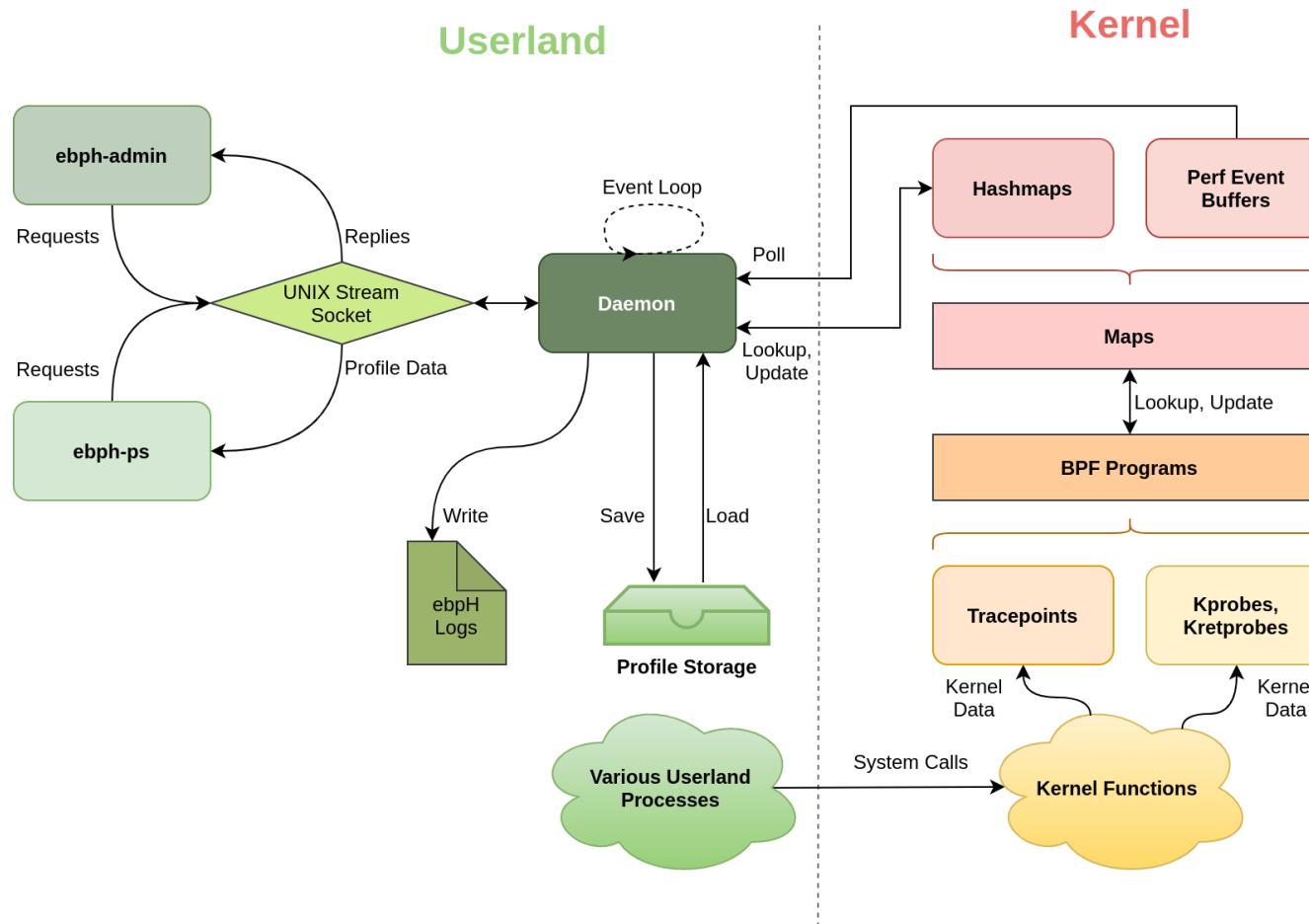


Figure 4: ebpH architecture in a nutshell.

- ◆ **How does ebpH overhead compare with pH?**
- ◆ **Benchmarks**
 - Imbench OS suite (micro)
 - ➔ *System call overhead*
 - ➔ *Process creation overhead*
 - ➔ *IPC overhead (signals, UDS, pipes)*
 - Kernel compilation benchmarks (micro)
 - ➔ *How does ebpH perform on real tasks?*
 - bpfbench (macro, ad-hoc)
 - ➔ *Real world system call overhead*
 - ➔ *Most frequent system calls in practice*

Performance Analysis

Table 2: Systems used for benchmarking tests.

System	Description	Specifications	
arch	Personal workstation	Kernel	5.5.10-arch1-1
		CPU	Intel i7-7700K (8) @ 4.500GHz
		GPU	NVIDIA GeForce GTX 1070
		RAM	16GB DDR4 3000MT/s
		Disk	1TB Samsung NVMe M.2 SSD
bronte	CCSL workstation	Kernel	5.3.0-42-generic
		CPU	AMD Ryzen 7 1700 (16) @ 3.000GHz
		GPU	AMD Radeon RX
		RAM	32GB DDR4 1200MT/s
		Disk	250GB Samsung SATA SSD 850
homeostasis	Mediawiki server	Kernel	5.3.0-42-generic
		CPU	Intel i7-3615QM (8) @ 2.300GHz
		GPU	Integrated
		RAM	16GB DDR3 1600MT/s
		Disk	500GB Crucial CT525MX3

Imbench Results

- ◆ **Short system calls**
 - `getppid(2)`: 614% overhead
→ *Almost no kernelspace runtime*
 - `stat(2)`: 65% overhead
→ *More significant kernelspace runtime*

- ◆ **Long system calls**
 - `select(2)`
 - As high as 99%
 - But as low as 2%

Figure 5: System call overheads.

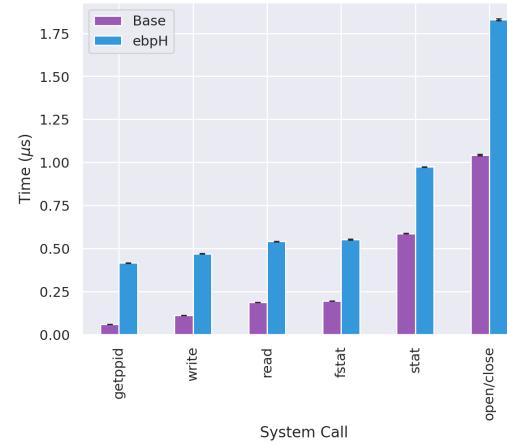
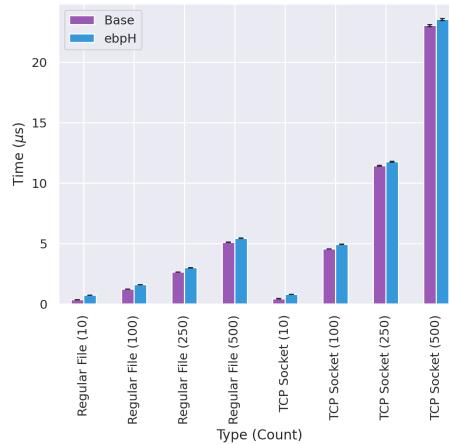


Figure 6: Various select(2) system call overheads.



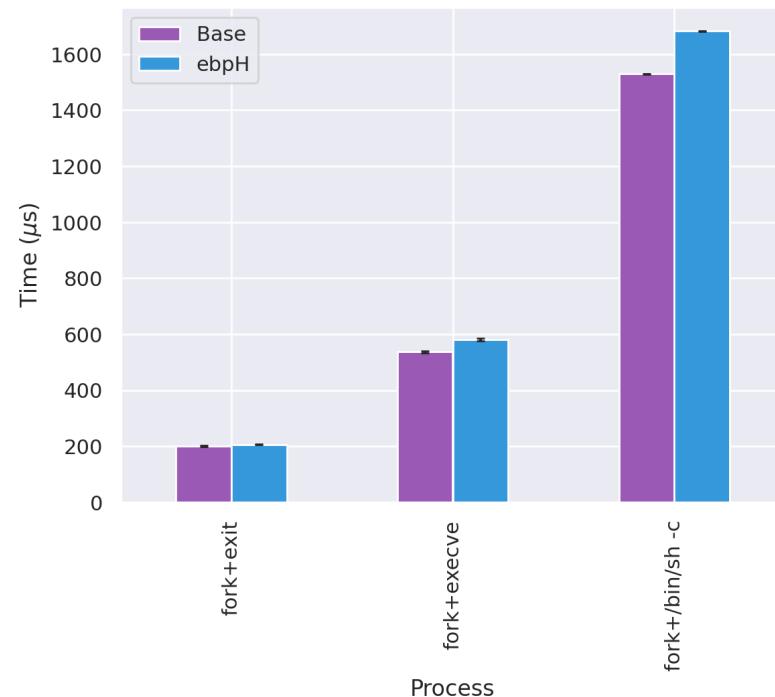
*Error bars show standard error.

Imbench Results

◆ Process creation

- fork+exit:
 - ➔ 2.7% overhead
- fork+execve:
 - ➔ 8.1% overhead
- fork+/bin/sh -c:
 - ➔ 10% overhead

Figure 7: Process creation latency results.
Least to most complex.



*Error bars show standard error.

Kernel Compilation Results

◆ Kernel compilation

- CPU-intensive task
- A lot of userspace time
- Still **many** system calls
 - ➔ Over **176 million**

◆ ebpH performs remarkably well here

- 10% kernelspace overhead
- 0.3% userspace overhead
- under 1% real overhead

Table 3: ebpH kernel compilation overheads.
Tests were run using 16 logical cores.

Category	T_{base} (s)	T_{ebpH} (s)	Diff. (s)	% Overhead
System	1525.412 (1.7603)	1687.833 (8.0621)	162.421667	10.647727
User	12333.737 (27.8529)	12370.957 (4.1244)	37.220000	0.301774
Elapsed	915.173 (3.9876)	924.032 (1.1194)	8.858333	0.967940

Table 4: Original pH kernel compilation overheads.

Time Category	Standard (s)	pH (s)	% Increase
user	728.92 (0.74)	733.09 (0.17)	0.57%
system	58.19 (0.80)	80.34 (0.17)	38.06%
elapsed	798.65 (0.87)	825.18 (1.75)	3.32%

*Standard deviations in parentheses.

- ◆ **Looked at top 20 system calls by count from three datasets**
 - arch (personal use)
 - bronte (idle)
 - homeostasis (production use)
- ◆ **Most frequent system calls have acceptable overhead**
 - Anywhere from about 5% to about 150%
- ◆ **Idle system reported significantly more overhead than the other two**
 - Lower overhead when it actually matters

- ◆ **ebpH imposes significant overhead on some system calls**
 - But this is not the whole story
 - ➔ *Longer system calls means less overhead*
 - ➔ *System call overhead ≠ overall impact*
- ◆ **Impact on most frequent system calls can be much lower in practice**
- ◆ **ebpH does very well on real tasks**
 - In some cases better than the original pH
 - Slowdown is mostly imperceptible in practice

◆ **bpf_signal**

- Real-time signals from kernelspace (**instantaneously**)
- SIGKILL, SIGSTOP, SIGCONT... you name it
- Linux 5.3

◆ **bpf_signal_thread**

- Like bpf_signal but target a specific thread
- Linux 5.5

◆ **bpf_override_return**

- Targeted **error injection**
- Whitelisted kernel functions only :(
- Linux 4.16

- ◆ **Add system call delays**
 - `bpf_signal` → send SIGSTOP and SIGCONT for delays
- ◆ **Add execve abortion**
 - `bpf_override_return` → target execve implementation

Table 1A: Adding response to ebpH.

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Future Work: Saving on Memory Overhead

- ◆ **Current map allocation is too granular**
 - One big map for profiles, one big map for processes
- ◆ **Solution: use new map types**
 - LRU_HASH → smaller map, discard least recently used entries
 - HASH_OF_MAPS → nested maps for lookahead pairs (sparse array)

Table 1B: Fixing ebpH's memory overhead.

System	Implementation	Portable	Production Safe	Low Mem. Overhead	Low Perf. Overhead	Detection	Response
pH	Kernel Patch	✗	✗	✓	✓	✓	✓
ebpH	eBPF + Userspace Daemon	✓	✓	✓	✓	✓	✓

- ◆ **Anomaly detection**
 - Add more sources of data?
 - No reason to stop at system calls
- ◆ **DDoS mitigation**
 - Cloudflare is doing this with eBPF/XDP
- ◆ **Increasing visibility of attacks / misuse**
 - ebpH does a bit of this
 - bcc tools are great for this
 - ➔ e.g. capable(8), eperm(8), setuids(8), execsnoop(8), etc.

- ◆ **Sandboxing?**

- Externally enforcing seccomp rules with eBPF?
- bpf_signal could do this easily

- ◆ **Name something you want to trace**

- eBPF can do it
- And it can do it **safely** and with **excellent performance**

- ◆ **ebpH is just the beginning**

- Uses a small fraction of eBPF's capabilities

Conclusion

◆ **ebpH:**

- is as fast as the original implementation
- supports most of the original functionality
- can be made even better, using new eBPF features

◆ **Future of ebpH?**

- Ecosystem of BPF programs
- All talking to each other, sharing information about diff. parts of system
- Beyond just system call tracing

◆ **Future of eBPF in OS security?**

- We are going to be seeing a lot more of this
- eBPF keeps getting better and better
- Replacing many in-kernel implementations with something safer, with less opportunity cost

Some Links

<https://github.com/iovisor/bcc>

<https://github.com/willfindlay/honors-thesis>

<https://github.com/willfindlay/ebph>

PRs welcome!

Thank you!