CS 5530

Database Systems Spring 2020

Relational Model

Keys

Library

Patrons

Name	CardNum
Joe	1
Ann	2
Ben	3
Dan	4

Inventory

Serial	ISBN
1001	978-0590353427
1002	978-0590353427
1003	978-0679732242
1004	978-0394823379
1005	978-0394823379
1006	978-0062278791

CheckedOut

Serial
1001
1004
1005
1006

Phones

CardNum	Phone
1	555-5555
2	666-6666
3	777-7777
4	888-888
4	999-9999

Titles

ISBN	Title	Author
978-0590353427	Harry Potter	Rowling
978-0679732242	The Sound and the Fury	Faulkner
978-0394823379	The Lorax	Seuss
978-0062278791	Profiles in Courage	Kennedy
978-0441172719	Dune	Herbert

Library

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Serial	ISBN
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1004	978-0394823379

CheckedOut

CardNum	Serial
1	1001
1	1004
4	1005

Phones

CardNum	Phone
1	555-5555
2	666-6666

How do we find all the books checked out by Joe?

Titles

ISBN	Title	Author
978-0590353427	Harry Potter	Rowling
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978-0394823379	The Lorax	Seuss
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Library

Patrons

Name	CardNum
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Inventory

Serial	BDIN
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1002	978-0590353427
1003	978-0679732242
1004	978-0394823379
1005	978-0394823379
1006	978-0662278791

CheckedOut

CardNum	Serial
1	1001
	1004
4	1005
4	1006

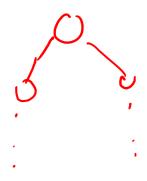
Phones

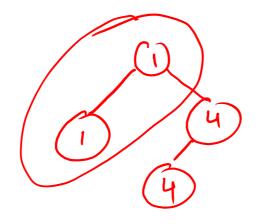
CardNum	Phone
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Titles

/	ISBN	Title	Author
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Don't Worry — It's Fast



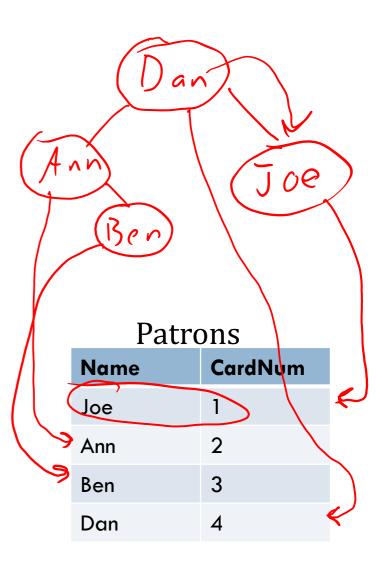


Inventory

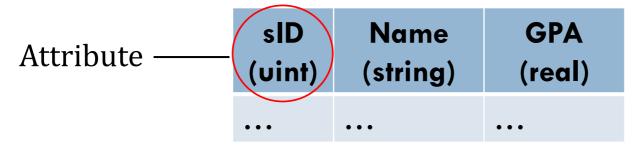
Serial	ISBN
1001	978-0590353427
1002	978-0590353427
1003	978-0679732242
1004	978-0394823379
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CheckedOut

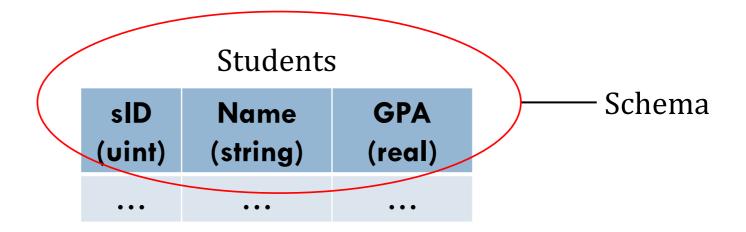
_	CardNum	Serial
	1	1001
	1	1004
	4	1005
	4	1006



•Attribute: a name and a type (column heading)



- •Schema: Table name + a set of attributes
 - Specifies the structure/rules of a table



•A schema does *not* specify any values

•Instance: the values in a table

• A set of *tuples*

	sID (uint)	Name (string)	GPA (real)
Instance—	1	Harry	3.5
	2	Hermione	4.0
	3	Ron	4.0
	4	Malfoy	3.9

•Instance: the values in a table

• A set of tuples

•Tuple: one row

	sID (uint)	Name (string)	GPA (real)	
{	1	Harry	3.5	} Tuple
	2	Hermione	4.0	
	3	Ron	4.0	
	4	Malfoy	3.9	

•Relation: a.k.a "table"

• Schema + instance

sID (uint)	Name (string)	GPA (real)
(01111)	(5111119)	(1001)
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

•Relation: a.k.a "table"

• A schema + instance

Relation1

sID (uint)	Name (string)	GPA (real)
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

Relation2

sID (vint)	Name (string)	GPA (real)
1	Harry	3.5
2	Hermione	4.0

- •Keys uniquely identify each tuple
 - Critical for the DBMS' underlying operations
 - Defines the meaning of the relation

Key	ı	
sID (uint)	Name (string)	GPA (real)
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

1---

	_ _	
Make	Model	•••
Toyota	F1	
Toyota	Camry	
Subaru	WRX	
Subaru	F1	

key

- •As a DB designer, you will define keys for each table
- •...but we need formal definitions before SQL

•Superkey:

•A set of attributes is a *superkey* if no two rows are allowed to have the same values in those columns (no duplicates)

Enrolled CID Grade SI

Student 5

SID Name GPA

(SIDICID) (SID, cID, Grade) (SID) (SID, ...)

•Key:

- •A set of fields is a *key* if:
 - it is a superkey
 - no proper subset of its fields is a superkey

•Key:

- •A set of fields is a *key* if:
 - it is a superkey
 - no proper subset of its fields is a superkey
- •(The empty set is not usually considered a superkey)

•What are the possible key(s)?

A ₁	A ₂	A_3
X	4	q
У	4	р
×	3	X

Attribute Set	SK	Key
$\{A_1\}$	No -	->
$\{A_2\}$	No -	
$\{A_3\}$	Yes	Yes
$\{A_1 A_2\}$	Yes	Yes
$\{A_1 A_3\}$	Yes	√ 0
$\{A_2 A_3\}$	Yes	
$\{A_1 A_2 A_3\}$	Yes	✓

•What are the possible key(s)?

A ₁	A ₂	A_3
X	4	q
У	4	р
X	3	X

Attribute Set	SK	Key
{A ₁ }	No	No
$\{A_2\}$	No	No
$\{A_3\}$	Yes	Yes
$\{A_1 A_2\}$	Yes	Yes
$\{A_1 A_3\}$	Yes	No
$\{A_2 A_3\}$	Yes	No
$\{A_1 A_2 A_3\}$	Yes	No

- •DB design works the other way around
 - Start with keys, not data
 - Keys define what data is valid

•DB designer's job: problem description → schemas

•What should be the key for the Enrolled schema?

Enrolled

sID	cID	Grade

- •sID = student ID
- •cID = course ID
- •No instance data given use human understanding of relation

•What should be the key for the Enrolled schema?

Enrolled

sID	clD	Grade

•What is the key for the Enrolled schema? Classes

cID Name when

• {sID, cID}

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

Enrolled

sID	cID	Grade	
1	C\$3500	В	
2	CS3500	A+	
4	CS3810	B-	
3	CS4400	A-	
2	CS6016	A+	

Primary/Candidate Key

DBA specifies one key as the primary key

•If a schema has more than one key, they are called candidate keys

Candidate Keys

•We can specify additional uniqueness constraints (candidate keys)



	sID	clD	Grade
--	-----	-----	-------

PK(sID, cID)
UQ(sID, Grade)

Candidate Keys

•We can specify additional uniqueness constraints (candidate keys)

Enrolled sID cID Grade

```
PK(sID, cID)
UQ(sID, Grade)
```

•"A student can not take a course twice, and can not receive the same grade twice"

•Provide the primary key (PK) plus any additional candidate keys (UQ) for the Enrolled table:



- 1. "A student can not take the same course twice, and no two students can receive the same grade in the same course"
- 2. "A student can earn multiple grades in a class"
- 3. "A student can only take one course and receive a single grade for that course, and courses have maximum enrollment of one"

• "A student can not take the same course twice, and no two students can receive the same grade in the same course"

"A student can retake a course for a different grade"

sID cID Grade

• "A student can only take one course and receive a single grade for that course, and courses have a maximum enrollment of one"

$$pk(sI0)$$
 $pk(sI0,cI0)$
 $uo(cI0)$

sID cID Grade

Primary vs. Candidate Keys

PK(sID, cID)

Or

UQ(sID, cID)

OR

UQ(sID, Grade)

PK(sID, Grade)

•"A student can not take a course twice, and can not receive the same grade in different courses"

Primary vs. Candidate Keys

```
PK(sID, cID)

Or

UQ(sID, cID)

PK(sID, Grade)

PK(sID, Grade)
```

•"A student can not take a course twice, and can not receive the same grade in different courses"

•PK should be small, ideally integer type(s)

•Keys relate tuples from different tables

Students			Enrolled		
sID	Name	GPA		sID	cID
1	Harry	3.5			
(2)-	Hermione	4.0		1	CS3500
3	Ron	4.0	→	2	CS3500
4	Malfoy	3.9		4	CS3810
	,			3	CS4400
				2	CS6016

Quiz

•What happens?

```
Thing* t1 = new Thing();
Thing* t2 = t1;
delete(t1);
cout << t2->x;
a) t2->value is printed
(b) crash
c) compiler error
```

d) no way to know what happens

Quiz

•What happens?

```
Thing* t1 = new Thing();
Thing* t2 = t1;
delete(t1);
cout << t2->x;
```

- a) t2->value is printed
- b) crash
- c) compiler error
- d) no way to know what happens

Quiz

•What happens?

```
Thing* t1 = new Thing();
Thing* t2 = t1;
delete(t1);
cout << t2->x;
```

- a) t2->value is printed
- b) crash what we *want* to happen (with useful message)
- c) compiler error
- d) no way to know what happens

Ponder...

Pointer analogy

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
3	CS4400	A-
2	CS6016	A+

Ponder...

•What happens if I delete the student "Ron"?

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	В-
3	CS4400	A-
2	CS6016	A+

Ponder...

•What happens if I delete the student "Ron"?

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sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

Enrolled

sID	clD	Grade
4	CS3810	B-
3	CS4400	A-
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meaningless

•Referential Integrity:

- References between values should always be meaningful
- This is an *invariant* should hold true at all times

•Referential Integrity:

- References between values should always be meaningful
- This is an *invariant* should hold true at all times

- •C++ lets us violate this
- •Java enforces it by taking away control

•Option 1:

• Whenever we remove a record from **Students**, run another command to update all **Enrolled** with that sID

•Option 1:

- Whenever we remove a record from **Students**, run another command to update all **Enrolled** with that sID
- Invariant briefly broken
- What if more tables reference **Students**?

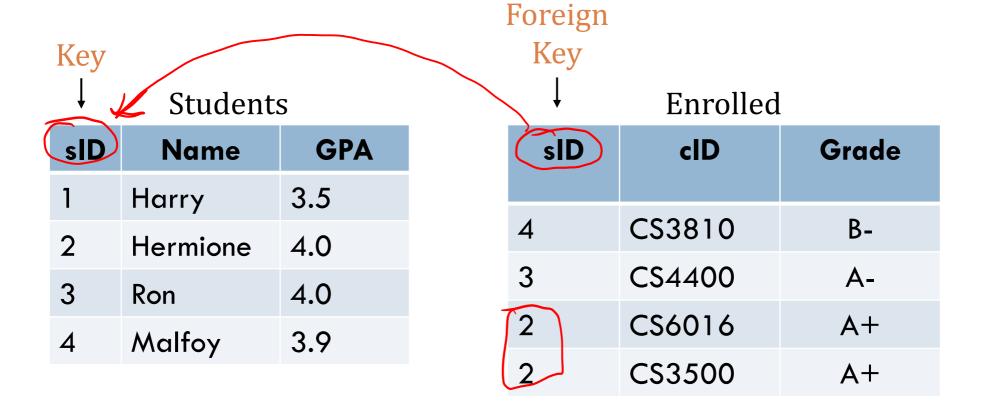
•Option 2:

- Let the DBMS update **Enrolled** automatically
- SQL provides some support

Foreign Key

•Foreign Key:

- Attribute whose values are a key in another table
- Think of it as a "pointer"



Foreign Key

•Foreign Key:

Not necessarily unique itself

Key Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

Foreign	
Key	
↓	Enrolled

sID	clD	Grade
4	CS3810	B-
3	CS4400	A-
2	CS6016	A+
2	CS3500	A+

SQL (ish)

```
CREATE TABLE Enrolled(
sID int,
cID char(6),
grade char(2),
PRIMARY KEY (sID, cID),
FOREIGN KEY (sID) REFERENCES Students(sID))
```

Enrolled

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	3.5
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
1	CS3500	Α
2	CS3500	В
4	CS3810	B-
3	CS4400	A-

Foreign Key Constraint

- •If the *referenced* key is modified, what should we do in the *referencing* table?
- •SQL gives a few options

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
3	CS4400	A-
2	CS6016	A+

Delete Record

•1. Delete corresponding record(s)



Patrons

Phones

Name	CardNum
Joe	1
Ann	2
Ben	3
Dan	4

CardNum	Phone
ī	555-5555
2	666-6666
3	777-7777
4	888-8888
4	999-9999

Delete Record

•1. Delete corresponding record(s)

Patrons

Phones

Name	CardNum
Joe	1
Ann	2
Ben	3
Dan	4

CardNum	Phone
1	555-5555
2	666-6666
3	777-7777
4	888-8888
4	999-9999

•If the referencing record has no meaning on its own

•2. Nullify foreign key

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
NULL	CS4400	A-
2	CS6016	A+

•2. Nullify foreign key

Students

sID	Name	GPA
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3	Ron	4.0
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sID	cID	Grade
4	CS3810	B-
NULL	CS4400	A-
2	CS6016	A+

- •If we want to keep the data, but "unlink" it
 - We can still analyze, e.g. average historic GPA

•2. Nullify foreign key

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	4.0
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
NULL	CS4400	A-
2	CS6016	A+

- •But! This is bad design
 - If {sID, cID} is primary key for Enrolled

•2. Nullify foreign key

Students

sID	Name	GPA
1	Harry	3.5
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3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
NULL	CS4400	A-
2	CS6016	A+

- •But! This is bad design
 - If {sID, cID} is primary key for Enrolled
- Add a unique ID to Enrolled table

Disallow

- •3. Disallow changes to referenced table
 - Try to delete Joe SQL reports error

Patrons

Name	CardNum
Joe	-======================================
Ann	2
Ben	3
Dan	4
steve	5

CheckedOut

CardNum	Serial
1	1001
1	1004
4	1005
4	1006

Disallow

- •3. Disallow changes to referenced table
 - Try to delete Joe SQL reports error

Patrons

Name	CardNum
Joe	-======================================
Ann	2
Ben	3
Dan	4

CheckedOut

CardNum	Serial
1	1001
1	1004
4	1005
4	1006

- •If we need to take some action first
 - Contact Joe, get his books back

Foreign Keys

- •What if we try to enroll a non-existent student?
 - SQL will reject it

Students

sID	Name	GPA
1	Harry	3.5
2	Hermione	3.5
3	Ron	4.0
4	Malfoy	3.9

sID	cID	Grade
4	CS3810	B-
3	CS4400	A-
2	CS6016	Α
5 (%)	CS4150	Е