

# 159.341 Programming Languages, Algorithms & Concurrency

C++ std::threads

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### Reminder

In the last lesson we looked at some of the basics of C++ threading.

- thread
- mutex
- timed\_mutex
- recursive\_mutex
- recursive\_timed\_mutex
- lock

## C++ std::threads - Mutexes

Mutexes are generally not intended to be used directly and more commonly utlised through lock\_guard or std::unique\_lock.

A lock\_guard is an RAII (Resource Allocation Is Initialisation) mechanism for taking ownership of a mutex for the duration of a scoped block.

A unique\_lock is a general-purpose mutex ownership wrapper that supports deferred locking, time-constrained attempts at locking, recursive locking, transfer of ownership and use with condition variables.

## C++ std::threads - lock\_guard

When a lock\_guard object is created, it will attempt to take ownership of the given mutex. When the lock\_guard leaves scope and is destructed, it will release the mutex.

```
std::mutex mtx:
void thread_function(int id) {
   std::this_thread::sleep_for(std::chrono::seconds(1));
   std::lock guard<std::mutex> lck(mtx);
   std::cout << "Thread: " << id << std::endl;</pre>
int main() {
   std::vector<std::thread> threads:
   for(int i = 0; i < 5; i++) {
      threads.push_back(std::thread(thread_function, i));
   for(std::thread &t : threads) {
     t.join();
```

## C++ std::threads - lock\_guard

One of the main advantages of using a lock\_guard is that they ensure mutexes are properly unlocked when an exception occurs.

If a lock\_guard is created inside a try block and an exception occurs. The mutex will be properly released when the exception-handling piece of code is executed and the lock\_guard leaves scope.

## C++ std::threads - lock\_guard

```
std::mutex mtx:
void print_maybe(int i) {
  int r = rand():
  if(r % 2 == 0) {
      std::cout << i << std::endl;
  } else {
      throw "Unlucky";
void thread function(int id) {
   trv {
      std::lock_guard<std::mutex> lock(mtx);
      print_maybe(id);
   } catch(const char *s) {
      // Handle exception
int main() {
   std::vector<std::thread> threads:
   for(int i = 0; i < 5; i++) {
      threads.push back(std::thread(thread function. i)):
   for(std::thread &t : threads) {
      t.join():
```

## C++ std::lock\_guard

There is another constructor for lock\_guard

lock\_guard(mutex\_type &m, std::adopt\_lock\_t t)

This constructor will create a lock\_guard and adopt the ownership of the mutex without trying to lock it.

Adopting the lock of a mutex should only ever be used if the thread already owns the mutex (behaviour is undefined otherwise).

## C++ - std::lock\_guard

The adopt\_lock constructor allows lock\_guards to be used in conjunction with std::lock. First the necessary mutexes are locked, then lock\_guards are constructed that adopt those locks.

```
std::mutex mtx1. mtx2. mtx3:
void thread function a() {
   std::lock(mtx1, mtx2, mtx3);
   std::lock_guard<std::mutex> lck1(mtx1, std::adopt_lock);
   std::lock_guard<std::mutex> lck2(mtx2, std::adopt_lock);
   std::lock_guard < std::mutex > lck3(mtx3, std::adopt_lock);
   std::cout << "Thread a" << std::endl:
void thread function b() {
   std::lock(mtx3, mtx2, mtx1);
   std::lock guard<std::mutex> lck1(mtx1. std::adopt lock);
   std::lock guard < std::mutex > lck2(mtx2. std::adopt lock):
   std::lock_guard < std::mutex > lck3(mtx3, std::adopt_lock);
   std::cout << "Thread b" << std::endl:
```

## C++ unique\_lock

C++ also has unique\_lock which supports *move semantics* to allow a unique\_lock to be *moved*.

This feature can be useful when a thread that currently holds a mutex wishes to pass ownership of that mutex to another thread.

The original thread can *move* the unique\_lock to the other thread and ensure that no other thread obtains ownership of it in between.

## C++ unique\_lock

#### Moving a unique\_lock to another thread:

```
std::mutex mtx1:
void thread function(int id, std::unique lock<std::mutex> &&lck) {
   std::this thread::sleep for(std::chrono::seconds(1)):
  std::cout << "Thread " << id << " running." << std::endl;
  std::this_thread::sleep_for(std::chrono::seconds(1));
int main() {
  // Create Threads
  std::vector<std::thread> threads:
  for(int i = 0; i < 5; i++) {
     // Get unique lock
      std::unique_lock<std::mutex> lck(mtx1):
     // Create Thread (move unique lock to thread)
      std::cout << "Lanching thread " << i << "." << std::endl:
      threads.push back(std::thread(thread function, i, std::move(lck)));
  // Join Threads
  for(std::thread &t : threads) {
     t.join();
```

The C++ threading API (before C++20) does not support semaphores.

Instead a condition\_variable can be used to block a thread until *notified* by another thread.

While a thread is waiting on a condition\_variable it will be blocked and not scheduled on the CPU.

condition\_variable provides the following functions:

- wait Wait until notified
- wait\_for Wait until notified or timeout
- wait\_until Wait until notified or time point is reached
- notify\_one Notify one thread waiting on condition variable.
- notify\_all Notify all threads waiting on condition variable.

#### Example:

```
std::mutex mtx:
std::condition_variable cond;
bool ready = false:
void thread_function(int id) {
   std::unique_lock<std::mutex> lock(mtx);
   std::cout << "Thread " << id << " waiting." << std::endl;
   while (!ready) cond.wait(lock):
   std::cout << "Thread " << id << " completed." << std::endl;
int main() {
   std::vector<std::thread> threads:
   for(int i = 0: i < 5: ++i) {
      threads.push_back(std::thread(thread_function, i));
   1
   std::this thread::sleep for (std::chrono::seconds(3)):
   readv = true:
   cond.notify_all():
   for(std::thread &t : threads) {
      t.ioin():
```

In the following thread function, the thread obtains a lock on the mutex mtx which will be held until the call to cond.wait(lock).

When the thread calls wait it will release the lock. When the thread is *notified* it will wake up and obtain a lock on mtx before continuing execution.

```
std::mutex mtx;
std::condition_variable cond;
bool ready = false;

void thread_function(int id) {
    std::unique_lock<std::mutex> lock(mtx);
    std::cout << "Thread " << id << " waiting." << std::endl;
    while(!ready) cond.wait(lock);
    ...
}</pre>
```

Note the following line is used when waiting on a condition variable.

```
while(!ready) cond.wait(lock);
```

In general, threads are only woken up when another thread calls either notify\_one or notify\_all. However, some implementations may produce spurious calls to wake up threads.

It is the programmer's responsibility to ensure the correct conditions to resume execution have been met.

To continue our Game of Life example we should look at how to implement our *phase-parallel* simulator.

We will need to explicitly create threads that will each update one section of the grid.

After each new generation is computed, the threads must synchronise with each other before continuing onto the next generation.

#### First a sequential implementation

```
unsigned char *buffer[2]:
buffer[0] = new unsigned char[N*N];
buffer[1] = new unsigned char[N*N]:
int r = 0: // Read/Write Index
int w = 1:
for(int k = 0; k < N*N; ++k) { // Initialise
  buffer[r][k] = rand() % 2;
for (int ix = 0: ix < N: ++ix) { // For each Column
       // Count neighbours
        int count = count neighbours(buffer[r], ix, iv, N);
       // Update
       buffer[w][iv*N + ix] = gol update(buffer[r][iv*N+ix], count);
  // Swap buffers
  r = !r:
  w = 1w:
```

#### Game of Life Functions:

```
// Count neighbours
int count neighbours (unsigned char *buffer, int ix, int iv, int N) {
   // Neighbouring Indexes
   int vm1 = (iv == 0) ? N-1 : iv-1;
   int vp1 = (iv == N-1) ? 0 : iv+1;
   int xm1 = (ix == 0) ? N-1 : ix-1;
   int xp1 = (ix == N-1) ? 0 : ix+1:
   // Count neighbours
   return buffer[ym1 * N + xm1] + buffer[ym1 * N + ix] + buffer[vm1 * N + xp1] +
          buffer[ iy * N + xm1] +
                                                      buffer[ iy * N + xp1] +
          buffer[vp1 * N + xm1] + buffer[vp1 * N + ix] + buffer[vp1 * N + xp1]:
// Came of Life Rules
unsigned char gol_update(unsigned char status, int count) {
   if ((status == 0) kk (count == 3)) {
      return 1:
   } else if((status == 1) && (count == 2 || count == 3)) {
      return 1:
   } else {
      return 0:
```

Creating threads to update different sections of the grid is relatively straightforward.

There is still the question of how many threads to use, which will we avoid by making the user decide.

```
std::thread *threads = new std::thread[T];
for(int t = 0; t < T; ++t) {
    threads[t] = std::thread(gol, ...);
}</pre>
```

```
// Thread Function
void gol(unsigned char *buffer[2], int N. int G. int id, int num threads) {
 int r = 0: // Read/Write index
 int w = 1:
 int start = ( id * N) / num_threads; // Start Row
  int end = ((id+1) * N) / num threads; // End Row
  for (int ig = 0: ig < G: ++ig) {
                                  // For each Generation
     for(int iv = start; iv < end; ++iv) { // For each Row (for this thread)
        for(int ix = 0: ix < N: ++ix) { // For each Column
           // Count neighbours
           int count = count_neighbours(buffer[r], ix, iy, N);
           // Update
           buffer[w][iv*N + ix] = gol update(buffer[r][iv*N+ix], count);
     // Barrier
     barrier(ig+1, num threads):
     // Swap buffers
     r = !r:
     w = !w:
```

To ensure that all the threads have updated their sections of the grid before continuing onto the next generation, we have used the following function:

barrier(ig+1, num\_threads)

Barriers are only defined in the C++20 standard and may not be available in many compilers.

We will need to implement our own barrier function.

```
std::condition_variable barrier_condition;
std::mutex barrier mutex:
int barrier count = 0:
int barrier generation = 0:
void barrier (int g, int num threads) {
  // Lock Muter
   std::unique_lock<std::mutex> lock(barrier_mutex);
  // Increment barrier counter
  ++barrier count:
  if(barrier_count == num_threads) {
      // Update barrier generation
      barrier_generation = g;
     // Reset barrier count
      barrier count = 0:
     // Notify all waiting threads
      barrier condition.notify all():
  } else {
     // Wait
      barrier_condition.wait(lock, [=]() {
        return barrier_generation == g;
     }):
```

# **Summary**

- lock\_guard
- unique\_lock
- condition\_variable
- Game of Life implementation