

159.341 Programming Languages, Algorithms & Concurrency

Lambda Calculus (Part 2)

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λ Calculus Reminder

So far we have covered the basics of λ -calculus:

- Abstraction and the basis of λ calculus
- Introduction to λ calculus syntax
- Normal order β reduction
- Syntax for function definitions
- α conversion for removing name-clashes
- η reduction for simplifying expressions

Developing a Functional Language

Lambda calculus is the ***model of computation*** that underlies functional languages.

To see how this can actually be developed into a useful functional language, we will look at the following features.

- Conditionals
- Integers
- Arithmetic operators

Conditionals

We will see how additional layers can be developed to create some higher-level (and easier to read) notation.

We will construct conditional expressions with truth values `true` and `false` and develop boolean operations `and`, `not` and `or`.

First we need to look at some functions that will help us along the way `select_first` and `select_second`.

Conditionals

Consider the function:

```
def select_first =  $\lambda$ first. $\lambda$ second.first
```

When this is applied to two expressions, `select_first` will return the first.

For example:

```
((select_first exp1) exp2) ==  
(( $\lambda$ first. $\lambda$ second.first exp1) exp2) =>  
( $\lambda$ second.exp1 exp2) =>  
exp1
```

Conditionals

Consider the function:

```
def select_second =  $\lambda$ first. $\lambda$ second.second
```

When this is applied to two expressions, `select_second` will return the second.

For example:

```
((select_second exp1) exp2) ==  
(( $\lambda$ first. $\lambda$ second.second exp1) exp2) =>  
( $\lambda$ second.second exp2) =>  
exp2
```

Conditionals

We can now use these functions to build up conditional expressions similar to the ternary expressions used in C:

```
condition ? expression1 : expression2
```

If the condition evaluates to true, then the first expression is selected and if the condition is false then the second expression is selected.

Conditionals

A similar expression can be constructed in λ calculus.

def cond = $\lambda e1.\lambda e2.\lambda c.((c\ e1)\ e2)$

Consider this function cond applied to two arbitrary expressions exp1 and exp2:

```
((cond exp1) exp2) ==  
(( $\lambda e1.\lambda e2.\lambda c.((c\ e1)\ e2)$  exp1) exp2) =>  
( $\lambda e2.\lambda c.((c\ exp1)\ e2)$  exp2) =>  
 $\lambda c.((c\ exp1)\ exp2)$ 
```


Conditionals

If this function is applied to `select_first`:

```
λc.((c exp1) exp2) select_first =>  
((select_first exp1) exp2) => ... =>  
exp1
```

or `select_second`:

```
(λc.((c exp1) exp2) select_second) =>  
((select_second exp1) exp2) => ... =>  
exp2
```

Conditionals

We can use the `cond` function to implement conditionals and the following definitions:

```
def cond =  $\lambda e1.\lambda e2.\lambda c.((c\ e1)\ e2)$   
def true = select_first =  $\lambda first.\lambda second.first$   
def false = select_second =  $\lambda first.\lambda second.second$ 
```

(Probably not the representation of true/false that you expected)

Conditionals

The NOT operator is a unary operator of the form NOT operand. It should also return the opposite of the boolean argument given to it - not true = false and not false = true.

Using our previous definitions, this function can be described by the following definition:

```
def not =  $\lambda x.((\text{cond false}) \text{ true}) \ x)$ 
```

Conditionals

This can be simplified:

```
((cond false) true) x) ==  
(((λe1.λe2.λc.((c e1) e2) false) true) x) =>  
((λe2.λc.((c false) e2) true) x) =>  
(λc.((c false) true) x) =>  
((x false) true)
```

Thus we use:

```
def not = λx.((x false) true)
```

Conditionals

For example (not true):

```
(not true) ==  
(λx.((x false) true) true) =>  
((true false) true) ==  
((λfirst.λsecond.first false) true) =>  
(λsecond.false true) =>  
false
```

Conditionals

Work through the similar evaluation of `(not false)`.

Conditionals

Work through the similar evaluation of `(not false)`.

```
(not false) ==  
(λx.((x false) true) false) =>  
((false false) true) ==  
((λfirst.λsecond.second false) true) =>  
(λsecond.second true) =>  
true
```

Conditionals

The `or` operator can also be constructed by using selectors. This function should result in `true` if either of the arguments are `true`.

If the first operand is `true` then select `true`, otherwise select the second operand.

```
def or =  $\lambda x. \lambda y. (((\text{cond } \text{true}) \text{ } y) \text{ } x)$ 
```

Like the operator `not`, this can be simplified.

Conditionals

```
((cond true) y) x) ==  
(((λe1.λe2.λc.((c e1) e2) true) y) x) =>  
((λe2.λc.((c true) e2) y) x) =>  
(λc.((c true) y) x) =>  
((x true) y)
```

Thus we have:

```
def or = λx.λy.((x true) y)
```

Conditionals

Evaluation of `((or true) false)`:

```
((or true) false) ==  
((λx.λy.((x true) y) true) false)  
(λy.((true true) y) false)  
((true true) false)  
((λfirst.λsecond.first true) false)  
(λsecond.true false)  
true
```

Conditionals

Try to write (and simplify) the and operator:

Conditionals

Try to write (and simplify) the and operator:

```
def and =  $\lambda x.\lambda y.(((\text{cond } y) \text{ false}) x)$ 
```

```
(((cond y) false) x) ==  
((( $\lambda e1.\lambda e2.\lambda c.((c \text{ e1}) e2) y$ ) false) x) =>  
(( $\lambda e2.\lambda c.((c y) e2) \text{ false}$ ) x) =>  
( $\lambda c.((c y) \text{ false}) x$ ) =>  
(x y) false)
```

Thus:

```
def and =  $\lambda x.\lambda y.(x y) \text{ false}$ 
```

Conditionals

Evaluation of `((and true) false)`:

```
((and true) false) ==  
((λx.λy.((x y) false) true) false) =>  
(λy.((true y) false) false) =>  
((true false) false) ==  
((λfirst.λsecond.first false) false) =>  
(λsecond.false false) =>  
false
```

Developing a Functional Language

We have looked at the simple rules of λ calculus and how we can define conditionals and logical operators.

There are a number of other features necessary to fully develop a functional language from λ calculus functions. Most important are integers, arithmetic operations and recursion.

There is also a significant amount of syntax simplification used for functional languages, but it should be unambiguous and able to be directly translated back into pure λ calculus.

Developing a Functional Language

One obvious requirement for a functional language is to support numbers and arithmetic.

Natural numbers (non-negative integers) can be defined a ***successors*** of ***zero***. 1 is the successor of 0, 2 is the successor of, the successor of 0, three is the successor of, the successor of, the successor of 0 (etc).

```
def one = (succ zero)
def two = (succ one)
def three = (succ two)
```

Developing a Functional Language

This allows natural numbers to be defined as that many successors of zero. Thus does require a function zero and succ. One way of representing these is to use the following definitions:

```
def zero = identity
def succ =  $\lambda n.\lambda s.((s\ false)\ n)$ 
```

This definition builds a pair function with `false` first and the number second.

Developing a Functional Language

For example:

```
one ==  
(succ zero) ==  
(λn.λs.((s false) n) zero) =>  
λs.((s false) zero)
```

```
two ==  
(succ one) ==  
(λn.λs.((s false) n) one) =>  
λs.((s false) one) ==  
λs.((s false) λs.((s false) zero))
```

Developing a Functional Language

Similarly a predecessor function `pred` can be constructed such that:

```
(pred one)    => ... => zero
(pred two)    => ... => one
(pred three) => ... => two
```

For our number representation, this can be constructed by removing a layer of nesting from a given number.

```
λs.((s false) number)
```

Developing a Functional Language

A first attempt at defining the predecessor function (we will call this `pred1`) could be:

```
def pred1 =  $\lambda n.(n \text{ select\_second})$ 
```

In the general application this would give:

```
(pred1  $\lambda s.((s \text{ false}) \text{ number})) ==$   
( $\lambda n.(n \text{ select\_second}) \lambda s.((s \text{ false}) \text{ number})) =>$   
( $\lambda s.((s \text{ false}) \text{ number}) \text{ select\_second}) =>$   
((select_second false) number) == ... =>  
number
```

Developing a Functional Language

What is the problem with this function definition?

Developing a Functional Language

What is the problem with this function definition?

What about zero?

```
(pred1 zero) ==  
(λn.(n select_second) zero) =>  
(zero select_second) ==  
(λx.x select_second) =>  
select_second ==  
false
```

Which does not represent a number.

Developing a Functional Language

Instead we could define that the predecessor of zero is zero and use the definition:

```
def pred =  $\lambda n.(((\text{cond zero}) (\text{pred1 } n)) (\text{iszero } n))$ 
```

Which can be simplified to:

```
def pred = (((iszero n) zero) (pred1 n))
```

```
def iszero =  $\lambda n.(n \text{ true})$ 
```

Summary

- Conditionals - true and false
- Logical Operators - not, and, or
- Defining Integers - succ, pred

Chapters 2 & 3

An Introduction to Functional Programming Through Lambda Calculus

Greg Michaelson