# 《并行计算》上机报告

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上机题目: MPI 并行编程实验

## 实验环境:

CPU: 8 个 Intel(R) Core(TM) i5-8250U

内存: 8G(交换区 8G)

操作系统: ubuntu 20.04 (内核版本为 5.4.0-72-generic)

软件平台: VSCode,

## 一、算法设计与分析:

## 题目一:

根据 PPT 上的内容,可以在求和近似积分的基础上,设计基于消息传递的并行求  $\pi$  的算法:

- (1) 非 0 号线程执行局部的求和, 并将其计算结果传递给 0 号线程。
- (2) 0号线程将这些结果进行累加,即得到π的近似值。

相比于 OpenMP 中共享存储的方法进行局部结果的累加,mpi 通过消息传递的方式需要考虑更多的开销,因此更慢一些。

#### 题目二:

根据教材上的内容,可以把PSRS分为以下8个阶段进行:

- (1) 均匀划分:将n个元素 A[1..n]均匀划分成p段,
- (2) 局部排序: pi 调用串行排序算法对 A[(i-1)n/p+1..in/p]排序
- (3) 选取样本: pi 从其有序子序列 A[(i-1)n/p+1..in/p]中选取 p 个样本元素
- (4) 样本排序: 用一台处理器对 p^2 个样本元素进行串行排序
- (5) 选择主元: 用一台处理器从排好序的样本序列中选取 p-1 个主元,并播送给其他 pi
- (6) 主元划分: pi 按主元将有序段 A[(i-1)n/p+1..in/p]划分成 p 段
- (7) 全局交换: 各处理器将其有序段按段号交换到对应的处理器中
- (8) 归并排序: 各处理器对接收到的元素进行归并排序

分析以上过程可知, (1)(4)(5)(7)(8)可以在一个线程上执行, (2)(3)(6)可以多个线程并行执行。而不同步骤之间需要进行一到多或者多到一的消息传递。

如果不考虑通信的开销,时间复杂度的分析和 OpenMP 的类似:

- (1) 均匀划分需要 O(p)的复杂度。
- (2) 局部排序需要 O( n/p lg(n/p) )的复杂度。
- (3) 选取样本需要 O(p)的复杂度。
- (4) 样本排序需要 O(p^2lg p)的复杂度。
- (5) 选择主元需要 O(p)的复杂度。
- (6) 主元划分需要 O(p)的复杂度。

- (7) 全局交换需要 O(p^2)的复杂度。
- (8) 归并排序需要 O(n/p lg(n/p))的复杂度。

因此,不考虑通信开销,基于 mpi 的 PSRS 算法的复杂度为 O(n/p lg(n/p) + p^2lg p); 但是,通信开销会使这里的复杂度继续增加。

## 二、核心代码:

#### 题目一:

(1) 在 main 函数中,控制不同线程号的线程执行不同的操作。

```
// for ID != 0, calculate the partial sum over step
if(thread_ID != 0){
    calculate_over_step(thread_ID,thread_num-1);
}
// for ID = 0, aggregate the partial sums
else{
    double pi = aggregate(thread_num);
    std::cout << "pi is " << std::setprecision(8) << pi << std::endl;
}</pre>
```

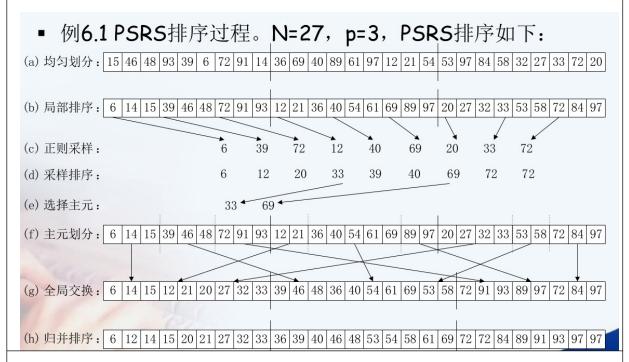
(2) 在非 0 号线程执行的 calculate\_over\_step()函数中,先进行循环的局部求和,再把求和的结果通过 MPI\_Send()发给 0 号线程。

```
// calculate the integral
// by getting sum of the interval over step
// @_thread_ID: current thread ID
// @_thread_num: total number of threads for computing integral
// send back value of the partial integral for aggregation
void calculate_over_step(int _thread_ID,int _thread_num){
    double x;
    double step = 1.0 / num_steps;
    double partial_sum = 0.0;
    // calculate the partial sum over step
    for(int i = _thread_ID;i < num_steps;i += _thread_num){
        x = (i + 0.5)*step;
        partial_sum += 4.0/(1.0 + x*x);
    }
    partial_sum *= step;
    // send the partial_sum to thread 0
    MPI_Send(&partial_sum,1, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD);
}</pre>
```

(3) 在 0 号线程执行的 aggregate()函数中,通过 MPI\_Recv() 接收每个非 0 线程的求和的结果,并将这些局部的结果进行累加。

## 题目二:

PSRS 算法的例子如下所示



由此可知, (1)均匀划分、(4)采样排序、(5)选择主元、(7)全局交换、(8)归并排序都是 在同一个线程里完成的。而(2)局部排序、(3)正则采样、(6)主元划分都是多线程并行完成 的。因此,我令 0 号线程执行 serial()串行函数,令非 0 号线程执行 parallel()并行函数, 完成以上操作;而核心代码如下:

一. 0号线程读取输入。

```
// get the input from the file
int elem_num;
int* array = get_input(elem_num);
// show the input array
std::cout << "original array is\n";
for(int i = 0; i < elem_num;++i){
    std::cout << array[i] << " ";
}
std::cout << "\n";</pre>
```

二. 0号线程对输入元素进

三. 每个非0号线程都得到得到它们对应的划分,并对它们进行排序。

四. 每一个非0号线程对划分进行采样, 并传递给0号线程。

```
175
176
          // (3) normal sampling
178
          // get thread num samples in a partition
179
          int gap = (int)floor( array size / thread num );
          int* local samples = new int[ thread num];
          int index = 0;
          for(int i = 0;i < _thread_num;++i){</pre>
183
              local_samples[i] = local_array[index];
184
              index += gap;
          // send the samples to thread 0 for next step
          MPI_Send(local_samples,_thread_num, MPI_INT, 0, 3, MPI_COMM_WORLD);
```

## 五.0号线程收集采样,并将采样进行排序

六. 0号线程选择主元,并将主元传递给每个非0号线程。

七.每一个非0号线程接收主元,并根据主元再次进行划分,用一个二维数组存放,并将这些划分传递给0号线程。

```
int *local pivots = new int[ thread num - 1];
          MPI Recv(local pivots, thread num - 1, MPI INT,
               0, 4, MPI COMM WORLD, MPI STATUS IGNORE);
          // replace list[i] means the elements to be sent to thread i
          std::array<std::vector<int>, MAX THREAD NUM> replace list;
          int pivot index = 0;
          for(int i = 0;i < array size;){</pre>
               if(local array[i] < local pivots[pivot index]){</pre>
203
                   replace list[pivot index].push back(local array[i]);
                  i++;
              else if(pivot index == thread num - 1){
                  replace list[pivot index - 1].push back(local array[i]);
                  i++;
210
              else{
211
                  pivot index++;
212
```

八. 0 号线程接收这些进一步的划分,并对他们进行归并排序,得到了最终有序的数组。

```
114
115
          // (7) global exchange & (8) merge sort
116
          std::array<std::vector<int>,MAX_THREAD_NUM> replace list;
119
          for (int thread index = 1; thread index <= thread num; ++ thread index){</pre>
120
              for(int part_index = 0; part_index < _thread_num; ++ part_index){</pre>
                  // get the partition in a thread
122
                  unsigned len;
                  MPI Recv(&len, 1, MPI UNSIGNED, thread index, 5, MPI COMM WORLD,
123
                  MPI STATUS IGNORE);
                  std::vector<int> partition(len, 0);
125
                  MPI_Recv(&partition[0], len, MPI_INT, thread_index, 6,
                  MPI COMM WORLD, MPI STATUS IGNORE);
126
                  int next part index = partition[len-1];
                  // tmp for merging
                  std::vector<int> tmp;
128
129
                  std::merge(replace_list[next_part_index].begin(), replace_list
                  [next part index].end(),
                       partition.begin(), --partition.end(), std::back_inserter(tmp));
130
131
                  replace list[next part index] = tmp;
132
```

#### 三、结果与分析:

#### 题目一:

我编译了程序之后,使用3个线程进行执行(实际上只有两个线程进行局部求和),得到了结果如下:

```
eddie@eddie-TM1701:~/Desktop/junior_2/parallel/LAB/lab2$ mpicxx -o par_pi par_pi.cpp
eddie@eddie-TM1701:~/Desktop/junior_2/parallel/LAB/lab2$ mpirun -np 3 ./par_pi
pi is 3.1415527
```

这里的结果精确到了小数点后4位。

#### 题目二:

使用 PPT 上的例子, 开 4 个线程执行, 得到测试结果如下:

```
eddie@eddie-TM1701:~/Desktop/junior_2/parallel/LAB/lab2$ mpicxx -o PSRS PSRS.cpp
eddie@eddie-TM1701:~/Desktop/junior_2/parallel/LAB/lab2$ mpirun -np 4 ./PSRS
original array is
15 46 48 93 39 6 72 91 14 36 69 40 89 61 97 12 21 54 53 97 84 58 32 27 33 72 20
after sorting, the array is
6 12 14 15 20 21 27 32 33 36 39 40 46 48 53 54 58 61 69 72 72 84 89 91 93 97 97
```

这说明可以正确排序。

# 四、备注(\* 可选):

有可能影响结论的因素:

- 1. 计算精度会导致估计π的结果并不是很准确
- 2. mpi 中的线程的消息传递和并行调度会产生开销,导致计算并不是很快。(也可能是输入的数据量较小,没有明显体现出并行编程的优越性出来)

## 总结:

本次实验, 我学习了使用 mpi 进行并行编程, 了解了基于消息传递的并行编程的思想, 并提高了自己的编码和调试能力。

附录(源代

算法源代码(C/C++/JAVA 描述)

我的本次实验的源码保存在了 https://github.com/fanweneddie/ 码) parallel\_computing/tree/master/lab2 中。 题目1: // par\_pi.c \* using MPI to calculate the approximate value of Pi \* by computing parallelism integral over step \* I have borrowed some ideas from myl7's github \* https://github.com/myl7/paracomp2021 \* PB18111679 fanweneddie #include <iostream> #include <iomanip> #include <mpi.h> static long num\_steps = 100000; void calculate\_over\_step(int \_thread\_ID,int \_thread\_num); double aggregate(int \_thread\_num); // calculate the integral // by getting sum of the interval over step // @\_thread\_ID: current thread ID // @\_thread\_num: total number of threads for computing integral // send back value of the partial integral for aggregation void calculate\_over\_step(int \_thread\_ID,int \_thread\_num){ double x; double step = 1.0 / num\_steps; double partial\_sum = 0.0; // calculate the partial sum over step for(int i = \_thread\_ID;i < num\_steps;i += \_thread\_num){ x = (i + 0.5)\*step;partial\_sum += 4.0/(1.0 + x\*x); partial\_sum \*= step; // send the partial\_sum to thread 0

```
// aggregate the partial sum computed by thread 1 \sim  _thread_num - 1
// @_thread_num: the number of threads
double aggregate(int _thread_num){
double pi = 0.0;
double partial_sum;
// receive the partial_sum calculated by other threads
// and aggregate them
for(int i = 0;i < _thread_num - 1;++i){
MPI_Recv(&partial_sum,1,MPI_DOUBLE,MPI_ANY_SOURCE,
0,MPI_COMM_WORLD,MPI_STATUS_IGNORE);
pi += partial_sum;
return pi;
int main(int argc, const char *const argv[]) {
MPI_Init(&argc, const_cast<char ***>(&argv));
// the thread ID of the current thread
int thread_ID;
MPI_Comm_rank(MPI_COMM_WORLD, &thread_ID);
// the total number of threads
int thread_num;
MPI_Comm_size(MPI_COMM_WORLD, &thread_num);
// for ID != 0, calculate the partial sum over step
if(thread_ID != 0){
calculate_over_step(thread_ID,thread_num-1);
// for ID = 0, aggregate the partial sums
double pi = aggregate(thread_num);
std::cout << "pi is " << std::setprecision(8) << pi << std::endl;
MPI_Finalize();
return 0;
```

MPI\_Send(&partial\_sum,1, MPI\_DOUBLE, 0, 0, MPI\_COMM\_WORLD);

```
题目二:
* use mpi to implement PSRS sorting algorithm
* the process of PSRS
* comes from our course slide
* I have borrowed some ideas from myl7's github
* https://github.com/myl7/paracomp2021
* PB18111679 fanweneddie
#include <fstream>
#include <iostream>
#include <array>
#include <vector>
#include <math.h>
#include <algorithm>
#include <mpi.h>
constexpr auto MAX_THREAD_NUM = 10;
void serial(int _thread_num);
void parallel(int _thread_ID,int _thread_num);
int* get_input();
// get the input int array from the input file
// allocate the array and return it
// @_elem_num: the number of elements in the array
// we need to revise it in the function
int* get_input(int& _elem_num){
std::ios::sync_with_stdio(false);
std::ifstream in("input.txt");
if (!in) {
std::cerr << "Getting input failed\n";
in >> _elem_num;
int *array = new int[_elem_num];
for (int i = 0; i < _elem_num; i++) {
in >> array[i];
in.close();
```

```
return array;
// serial working part
// for thread 0
// do the serial work in PSRS,
// that is, to do
// (1) uniform partition
// (4) sample sorting
// (5) pivot select
// (7) global exchange
// (8) merge sort
// @_thread_num: number of parallel threads
void serial(int _thread_num){
// get the input from the file
int elem_num;
int* array = get_input(elem_num);
// show the input array
std::cout << "original array is\n";
for(int i = 0; i < elem_num;++i){
std::cout << array[i] << " ";
std::cout << "\n";
// (1) uniform partition
// the index of start and end for each partition
int start_index,end_index;
// the gap between two neighbor partition
int gap = (int)ceil(elem_num / _thread_num );
for(int i = 1; i <= _thread_num;++i){
start_index = gap * (i - 1);
end_index = gap * i < elem_num ? gap * i : elem_num;
int partition_size = end_index - start_index;
// send the elements of each partition
// to the corresponding thread
MPI_Send(&partition_size,1,MPI_INT, i, 1, MPI_COMM_WORLD);
MPI_Send(array + start_index,
end_index - start_index, MPI_INT, i, 2, MPI_COMM_WORLD);
```

```
// receive the samples from those threads
int* samples = new int [_thread_num * _thread_num];
// the offset in samples for each reception
int offset = 0;
for(int i = 0;i < _thread_num;++i){
MPI_Recv(samples + offset, _thread_num,
MPI_INT,MPI_ANY_SOURCE, 3, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
offset += _thread_num;
// (4) sample sorting
std::sort(samples,samples + offset);
// (5) pivot select
int *pivots = new int[_thread_num - 1];
// get the pivots uniformly
// the step is _thread_num
for(int i = 0; i < _thread_num - 1;++i){
pivots[i] = samples[ (i + 1) * _thread_num ];
// send the pivots to thread 1 ~ _thread_num
for(int i = 1; i <= _thread_num;++i){
MPI_Send(pivots, thread_num - 1,MPI_INT, i, 4, MPI_COMM_WORLD);
// (7) global exchange & (8) merge sort
// receive the partitions sent by thread 1 ~_thread_num - 1
std::array<std::vector<int>,MAX_THREAD_NUM> replace_list;
for (int thread_index = 1; thread_index <= _thread_num; ++ thread_index){
for(int part_index = 0; part_index < _thread_num; ++ part_index){
// get the partition in a thread
unsigned len;
MPI_Recv(&len, 1, MPI_UNSIGNED, thread_index, 5, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
std::vector<int> partition(len, 0);
MPI_Recv(&partition[0], len, MPI_INT, thread_index, 6, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
int next_part_index = partition[len-1];
// tmp for merging
```

```
std::merge(replace_list[next_part_index].begin(), replace_list[next_part_index].end(),
partition.begin(), --partition.end(), std::back_inserter(tmp));
replace_list[next_part_index] = tmp;
// new array to store the sorted elements
std::vector<int> new_array;
// concat the replace_lists to get the new_array
for(int i = 0; i < _thread_num;++i){
new_array.insert(new_array.end(),replace_list[i].begin(),replace_list[i].end());
std::cout << "after sorting, the array is\n";
for (auto elem : new_array) {
std::cout << elem << " ";
std::cout << std::endl;
delete []array;
delete []samples;
delete []pivots;
// parallel working part
// for thread 1 ~ _thread_num - 1
// do the parallel work in PSRS
// that is, to do
// (2) local sort
// (3) uniform sampling
// (6) privot partition
void parallel(int _thread_ID,int _thread_num){
// receive the uniform partition from thread 0
int arrav size:
MPI_Recv(&array_size,1, MPI_INT, 0, 1, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
int* local_array = new int[array_size];
MPI_Recv(local_array, array_size, MPI_INT, 0, 2,
MPI_COMM_WORLD, MPI_STATUS_IGNORE);
// (2) local sort
std::sort(local_array,local_array + array_size);
```

std::vector<int> tmp;

```
// (3) normal sampling
// get _thread_num samples in a partition
int gap = (int)floor( array_size / _thread_num );
int* local_samples = new int[_thread_num];
int index = 0;
for(int i = 0;i < _thread_num;++i){
local_samples[i] = local_array[index];
index += gap;
// send the samples to thread_0 for next step
MPI_Send(local_samples,_thread_num, MPI_INT, 0, 3, MPI_COMM_WORLD);
// (6) privot partition
// receive the pivots sent from the first partition
int *local_pivots = new int[_thread_num - 1];
MPI_Recv(local_pivots,_thread_num - 1, MPI_INT,
0, 4,MPI_COMM_WORLD, MPI_STATUS_IGNORE);
// partition according to the pivots
// replace_list[i] means the elements to be sent to thread i
std::array<std::vector<int>, MAX_THREAD_NUM> replace_list;
int pivot_index = 0;
for(int i = 0;i < array_size;){
if(local_array[i] < local_pivots[pivot_index]){
replace_list[pivot_index].push_back(local_array[i]);
į++;
else if(pivot_index == _thread_num - 1){
replace_list[pivot_index - 1].push_back(local_array[i]);
j++:
else{
pivot_index++;
// send the result of partition to thread 0
for(int i = 0;i < _thread_num;++i){
int len = (unsigned)replace_list[i].size();
// marks the object of this partition.
```

```
replace_list[i].push_back(i);
len++;
MPI_Send(&len, 1, MPI_UNSIGNED, 0, 5, MPI_COMM_WORLD);
MPI_Send(&replace_list[i][0], len, MPI_INT, 0, 6, MPI_COMM_WORLD);
delete [] local_array;
delete [] local_samples;
delete [] local_pivots;
int main(int argc, const char *const argv[]) {
MPI_Init(&argc, const_cast<char ***>(&argv));
// the thread ID of the current thread
int thread_ID;
MPI_Comm_rank(MPI_COMM_WORLD, &thread_ID);
// the total number of threads
int thread_num;
MPI_Comm_size(MPI_COMM_WORLD, &thread_num);
thread_num = 4;
// for ID = 0, do the serial work
if(thread_ID == 0){
serial(thread_num - 1);
// for ID != 0, do parallel work
else{
parallel(thread_ID,thread_num - 1);
MPI_Finalize();
return 0;
}
```