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region and then type : Hash to hash your selection."

"Useful for verifying that there aren't mistypes.

"Select

```
$/d" \| tr -d '[:space:]' \| md5sum \| cut -c-6
1.2 readchar [a419b9]
inline char readchar() {
  static const size_t bufsize = 65536;
   static char buf[bufsize];
   static char *p = buf, *end = buf;
  if (p == end) end = buf +
         fread_unlocked(buf, 1, bufsize, stdin), p = buf;
   return *p++;
1.3 BigIntIO [d9afcb]
___int128 read() {
       _{\text{int}128 \ x} = 0, \ f = 1;
     char ch = getchar();
     while (ch < '0' || ch > '9') {
    if (ch = '-') f = -1;
          ch = getchar();
     while (ch >= '0' && ch <= '9') {
          x = x * 10 + ch - '0';
          ch = getchar();
     return x * f;
void print(__int128 x) {
     if (x < 0) {
putchar('-');
          x = -x:
     if (x > 9) print(x / 10);
putchar(x \% 10 + '0');
bool cmp(\underline{\phantom{a}}int128 x, \underline{\phantom{a}}int128 y) { return x > y; }
1.4 Black Magic [6547c5]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
      tree<int, null_type, std::less<int>, rb_tree_tag
, tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int , int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree set s:
   s.insert(71); s.insert(22);
   assert(*s.find_by_order
       (0) = 22; assert(*s.find_by_order(1) = 71);
  assert(s.order_of_key (22) == 0); assert(s.order_of_key(71) == 1);
  s.erase(22);
assert(*s.find_by_order
        (0) = 71; assert (s.order_of_key(71) = 0);
   // mergable heap
  heap a, b; a.join(b);
   // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
   std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0;
}
1.5 Pragma Optimization [eac636]
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse, sse2, sse3, sse4")
#pragma GCC target("popent, abm,mx, avx, arch=skylake")
```

_builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)

1.6 Bitset [84082f]

```
#include<bits/stdc++.h>
using namespace std;

int main () {
    bitset<4> bit;
    bit.all(); // all bit is true, ret tru;
    bit.any(); // any bit is true, ret true
    bit.none(); // all bit is false, ret true
    bit.count();
    bit.to_string('0', '1');//with parmeter
    bit.reset(); // set all to true
    bit.set(); // set all to false
    std::bitset<8> b3{0}, b4{42};
    std::hash<std::bitset<8> hash_fn8;
    hash_fn8(b3); hash_fn8(b4);
}
```

2 Graph

2.1 BCC Vertex* [740acb]

```
struct BCC { // 0-base
   int n, dft, nbcc;
   vector<int> low, dfn, bln, stk, is_ap, cir;
   vector<vector<int>>> G, bcc, nG;
   void make_bcc(int u) {
      bcc.emplace_back(1, u);
      for (; stk.back() != u; stk.pop_back())
  bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
      stk.pop\_back(), bln[u] = nbcc++;
   void dfs(int u, int f) {
      int child = 0;
     low [u] = dfn [u] = ++dft, stk.pb(u);
for (int v : G[u])
if (!dfn [v]) {
            dfs(v, u), +child;

low[u] = min(low[u],
                                           low[v]);
             \begin{array}{ll} & \text{if } (dfn [u] <= low [v]) \\ & \text{is\_ap} [u] = 1, \ bln [u] = nbcc; \\ & \text{make\_bcc}(v), \ bcc.back().pb(u); \\ \end{array} 
         else if (dfn[v] < dfn[u] & v != f)
     \begin{array}{l} low[u] = min(low[u], dfn[v]); \\ if \ (f = -1 \&\& child < 2) \ is\_ap[u] = 0; \\ if \ (f = -1 \&\& child = 0) \ make\_bcc(u); \end{array}
  BCC(int _n): n(_n), dft()
  G[u].pb(v), G[v].pb(u);
  void solve() {
  for (int i = 0; i < n; ++i)</pre>
         if (!dfn[i]) dfs(i, -1);
   void block_cut_tree() {
      cir.resize(nbcc);
      for (int i = 0; i < n; ++i)
         if (is_ap[i])
            bln[i] = nbcc++;
      cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)
    for (int j : bcc[i])</pre>
  if (is_ap[j])
    nG[i].pb(bln[j]), nG[bln[j]].pb(i);
} // up to 2 * n - 2 nodes!! bln[i] for id
```

2.2 Bridge* [4da29a]

```
struct ECC { // 0-base
  int n, dft, ecnt, necc;
  vector<int> low, dfn, bln, is_bridge, stk;
  vector<vector<pii>>> G;
  void dfs(int u, int f) {
    dfn[u] = low[u] = ++dft, stk.pb(u);
    for (auto [v, e] : G[u])
      if (!dfn[v])
        dfs(v, e), low[u] = min(low[u], low[v]);
      else if (e != f)
        low[u] = min(low[u], dfn[v]);
    if (low[u] = dfn[u]) {
      if (f != -1) is_bridge[f] = 1;
      for (; stk.back() != u; stk.pop_back())
```

```
b \ln \left[\, \mathrm{stk} \, . \, \mathrm{back} \, (\,) \, \right] \; = \; \mathrm{necc} \, ;
           bln[u] = necc++, stk.pop\_back();
       }
    ECC(int _n): n(_n), dft()
    , ecnt(), necc(), low(n), dfn(n), bln(n), G(n) {}
void add_edge(int u, int v) {
       G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
    void solve() {
       is_bridge.resize(ecnt);
for (int i = 0; i < n; ++i)
           if (!dfn[i]) dfs(i, -1);
 }; // ecc_id(i): bln[i]
 2.3 SCC* [4057dc]
 struct SCC { // 0-base
    int n, dft, nscc;
vector<int> low, dfn, bln, instack, stk;
     vector<vector<int>>> G;
    void dfs(int u) {
       low[u] = dfn[u] = ++dft;
       \begin{array}{l} instack \left[u\right] = 1, \ stk.pb(u); \\ for \ (int \ v \ : G[u]) \end{array}
           \begin{array}{l} \mbox{if } (!\,dfn\,[v]) \\ \mbox{dfs}(v)\,,\,\,low\,[u] \,=\, min(low\,[u]\,,\,\,low\,[v])\,; \\ \mbox{else} \mbox{if } (instack\,[v]\,\,\&\&\,\,dfn\,[v] \,<\,dfn\,[u]) \end{array}
        \begin{array}{c} low\left[u\right] = min(low\left[u\right], \ dfn\left[v\right]); \\ if \ (low\left[u\right] \Longrightarrow dfn\left[u\right]) \ \{ \end{array}
           for (; stk.back() != u; stk.pop_back())
              bln[stk
           \label{eq:continuous_continuous} \begin{array}{ll} \text{`.back()]} = \operatorname{nscc}, \ \operatorname{instack}\left[\operatorname{stk.back()}\right] = 0; \\ \operatorname{instack}\left[u\right] = 0, \ \operatorname{bln}\left[u\right] = \operatorname{nscc++}, \ \operatorname{stk.pop\_back()}; \end{array}
    void add_edge(int u, int v) {
       G[u].pb(v);
    void solve() {
  for (int i = 0; i < n; ++i)</pre>
           if (!dfn[i]) dfs(i);
 }; // scc_id(i): bln[i]
 2.4 2SAT* [f5630a]
struct SAT { // 0-base
    int n:
    vector<bool> istrue;
    SCC \ scc;
    SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
    int rv(int a) {
        return a >= n ? a - n : a + n;
    void add_clause(int a, int b) {
        scc.add\_edge(rv(a), b), scc.add\_edge(rv(b), a);
    bool solve() {
        scc.solve();
        for (int i = 0; i < n; ++i) {
           if (scc.bln[i] = scc.bln[i+n]) return false;
           istrue[i] = scc.bln[i] < scc.bln[i + n];
```

$2.5 \quad Minimum Mean Cycle* \ \tiny{\texttt{[3e5d2b]}}$

istrue[i + n] = !istrue[i];

return true;

};

```
for (int j = 1; j < n; ++j)
    if (dp[j][i] < INF &&
        ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
        ta = dp[L][i] - dp[j][i], tb = L - j;
    if (ta == 0) continue;
    if (a == -1 | | a * tb > ta * b) a = ta, b = tb;
}
if (a != -1) {
    ll g = __gcd(a, b);
    return pll(a / g, b / g);
}
return pll(-1LL, -1LL);
}
void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
}
};</pre>
```

2.6 Virtual Tree* [1b641b]

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
  if (top = -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
  if (p = st[top]) return st[++top] = u, void();
  while (top >= 1 \& dep[st[top - 1]] >= dep[p])
  vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
   vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
 vG[u].clear();
void solve(vector<int> &v) {
  sort (ALL(v),
  [&](int a, int b) { return dfn[a] < dfn[b]; });
for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do something
  reset (v[0]);
```

2.7 Maximum Clique Dyn* [d50aa9]

```
struct MaxClique { // fast when N \le 100
  void init(int _n) {
    n\,=\,\underline{}\,n;
    for (int i = 0; i < n; ++i) G[i].reset();
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  \begin{tabular}{ll} void $$ pre\_dfs(vector<int> \&r, int l, bitset< \gg mask) $$ \{ \end{tabular} \label{table}
    if (1 < 4) {
for (int i : r) d[i] = (G[i] & mask).count();
       sort (ALL(r)
            , [\&](int x, int y) \{ return d[x] > d[y]; \});
    vector < int > c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
       int k = 1;
       while ((cs[k] \& G[p]).any()) ++k;
        if (k > rgt) cs[++rgt + 1].reset(); \\
       cs[k][p] = 1;
       if (k < lft) r[tp++] = p;
    for (int k = lft; k \ll rgt; ++k)
       for (int p = cs[k]._Find_first
            (); p < N; p = cs[k]. \underline{\text{Find}}\underline{\text{next}}(p))
    r[tp] = p, c[tp] = k, ++tp;

dfs(r, c, l + 1, mask);
  void dfs (vector<
       int > &r, vector < int > &c, int 1, bitset < N > mask) {
```

```
while (!r.empty()) {
       int p = r.back();
       r.pop\_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;
       \operatorname{cur}\left[q++\right] = p;
       vector<int> nr;
for (int i : r) if (G[p][i]) nr.pb(i);
        if \quad (!\, nr.\, empty()) \quad pre\_dfs(nr, \ l \ , \ mask \ \& \ G[p]) \ ; \\
       else if (q > ans) ans = q, copy_n(cur, q, sol);
       c.pop_back(), --q;
    }
  int solve() {
     vector < int > r(n);
     ans = q = 0, iota (ALL(r), 0);
     pre_dfs(r, 0, bitset < N>(string(n, '1')));
     return ans;
};
```

2.8 Minimum Steiner Tree* [62d6fb]

```
struct SteinerTree { // 0-base
  \begin{array}{l} int \ n, \ dst \left[N\right]\left[N\right], \ dp\left[1 << T\right]\left[N\right], \ tdst \left[N\right]; \\ int \ vcst \left[N\right]; \ // \ the \ cost \ of \ vertexs \end{array}
  n = _n; for (int i = 0; i < n; ++i) {
        fill_n (dst[i], n, INF);
        dst[i][i] = vcst[i] = 0;
   void chmin(int &x, int val) {
    x = \min(x, val);
   void add_edge(int ui, int vi, int wi) {
     chmin(dst[ui][vi], wi);
   void shortest_path() {
     for (int k = 0; k < n; ++k)
for (int i = 0; i < n; ++i)
for (int j = 0; j < n; ++j)
             chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1;
     for (int i = 0; i \le full; ++i)
        fill_n (dp[i], n, INF);
     copy_n(vest, n, dp[0]);
for (int msk = 1; msk <= full; ++msk) {
  if (!(msk & (msk - 1))) {</pre>
          int who = \__lg(msk);
          for (int i = 0; i < n; ++i)
             \mathrm{dp}\,[\,\mathrm{msk}
                   [i] = vcst[ter[who]] + dst[ter[who]][i];
        for (int i = 0; i < n; ++i)
          for (int sub = (
             msk - 1) & msk; sub; sub = (sub - 1) & msk) chmin(dp[msk][i],
        for (int j = 0; j < n; ++j)

chmin(tdst[i], dp[msk][j] + dst[j][i]);
        copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
;'// O(V 3^T + V^2 2^T)
```

2.9 Dominator Tree* [2b8b32]

```
void dfs(int u) {
    id[dfn[u] = ++Time] = u;
for (auto v : G[u])
      if(!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  int find(int y, int x) {
    if (y \le x) return y;
    int tmp = find(pa[y], x);
    if (semi[best[y]] > semi[best[pa[y]]])
      best[y] = best[pa[y]];
    return pa[y] = tmp;
  void tarjan(int root) {
    Time = 0;
    for (int i = 1; i <= n; ++i) {
  dfn[i] = idom[i] = 0;
      tree[i].clear();
      best[i] = semi[i] = i;
    dfs(root);
for (int i = Time; i > 1; --i) {
      int u = id[i];
for (auto v : rG[u])
        if (v = dfn[v]) 
           tree[semi[i]].pb(i);
for (auto v : tree[pa[i]]) {
        find(v, pa[i]);
idom[v] =
           semi[best[v]] = pa[i] ? pa[i] : best[v];
      tree [pa[i]]. clear();
    for (int i = 2; i \leftarrow Time; ++i)
      if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
      tree [id [idom[i]]].pb(id[i]);
 }
2.10 Four Cycle [584a52]
```

```
int main() {
         cin.tie(nullptr)->sync_with_stdio(false);
         cin >> n >> m;
         for (int i = 1; i \le m; i++) {
                int u, v;
cin >> u >> v;
                E[u]. push_back(v);
                E[v]. push_back(u);
                 deg[u]++, deg[v]++;
         for (int u = 1; u \le n; u++)
                for (int u = 1, u
         for (int a = 1; a \le n; a++) {
                 for (int b : E1[a])
                          for (int c : E[b])
                                   if (deg[a] < deg[c]
                                                         | | (deg[a] = deg[c] \& a <= c)) continue;
                                  total += cnt[c]++;
                 for (int b : E1[a])
                          for (int c : E[b]) cnt[c] = 0;
        cout << total << '\n';
         return 0;
```

2.11 Minimum Clique Cover* [879472]

```
void init(int _n) {
    n = _n, fill_n(dp, 1 << n, 0);
  fill_n(E, n, 0), fill_n(co, 1 << n, 0);
 int solve() {
  for (int i = 0; i < n; ++i)
```

```
co[1 << i] = E[i] | (1 << i);
        co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
         for (int i = 1; i < (1 << n); ++i) {
            \begin{array}{lll} & \text{int } t = i \& -i; \\ & \text{dp}[i] = -\text{dp}[i \land t]; \\ & \text{co}[i] = \text{co}[i \land t] \& \text{co}[t]; \end{array}
         for (int i = 0; i < (1 << n); ++i)
             co[i] = (co[i] \& i) = i;
        fwt(co, 1 << n, 1);
for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)
    sum += (dp[i] *= co[i]);</pre>
             if (sum) return ans;
         return n;
};
```

NumberofMaximalClique* [11fa26] 2.12

```
struct BronKerbosch { // 1-base
  void init(int _n) {
     n = _n;
for (int i = 1; i \le n; ++i)
        for (int j = 1; j \le n; ++j) g[i][j] = 0;
   void add_edge(int u, int v) {
     g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
     if (sn = 0 \&\& nn = 0) ++S;
     int u = some[d][0];
for (int i = 0; i < sn; ++i) {
        \underset{\quad \text{int }}{\text{int }} v = some [d][i];
        if (g[u][v]) continue;
        int tsn = 0, tnn = 0
        copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
        for (int j = 0; j < sn; ++j)
if (g[v][some[d][j]])
            some [d + 1][tsn++] = some [d][j];
        for (int j = 0; j < nn; ++j)
  if (g[v][none[d][j]])</pre>
             none[d + 1][tnn++] = none[d][j];
        dfs(d + 1, an + 1, tsn, tnn);

some[d][i] = 0, none[d][nn++] = v;
  int solve() {
  iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
     return S;
};
```

3 Data Structure

3.1Discrete Trick

```
vector<int> val;
// build
sort (ALL
    (val)), val.resize(unique(ALL(val)) - val.begin());
upper_bound(ALL(val), x) - val.begin();
// \max idx \le x
upper_bound(ALL(val), x) - val.begin();
 / \max idx <
lower_bound(ALL(val), x) - val.begin();
3.2 BIT kth* [e39485]
```

int bit [N + 1]; // $N = 2^k$ int query_kth(int k) {

```
int res = 0;
     for (int i = N >> 1; i >= 1; i >>= 1)
         if (bit[res + i] < k)
             k \rightarrow bit[res += i];
    return res + 1;
}
```

3.3 Interval Container* [c54d29]

```
/* Add and
     remove intervals from a set of disjoint intervals.
* Will merge the added interval with
      any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>>::
    iterator addInterval(set<pii> is, int L, int R) {
  if (L = R) return is .end();
  {\color{red} \textbf{auto}} \ it = is.lower\_bound(\{L,\ R\})\,, \ before = it\,;
  while (it != is.end() && it->X<= R) {
    R = \max(R, it ->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y>= L) {
    L = min(L, it ->X);

R = max(R, it ->Y);
    is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii> is, int L, int R) {
 if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it ->Y;
  if (it->X == L) is .erase(it);
  else (int&)it-\hat{Y} = L;
  if (R != r2) is .emplace(R, r2);
```

3.4 Leftist Tree [e91538]

```
struct node {
   ll v, data, sz, sum;
node *1, *r;
   node(ll k)
      : v(0), data(k), sz(1), l(0), r(0), sum(k) \{ \}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
   if (a->data < b >data) gwar(a b);
   if (a->data < b->data) swap(a, b);
   a \rightarrow r = merge(a \rightarrow r, b);
   if (V(a->r) > V(a->l)) swap(a->r, a->l);
   a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
   a->sum = sum(a->l) + sum(a->r) + a->data;
   return a;
void pop(node *&o) {
  node *tmp = o;
   o = merge(o->l, o->r);
   delete tmp;
```

3.5 Heavy light Decomposition* [b004ae]

```
vector<int>G[N];
  void init(int _n) {
    n = _n;
for (int i = 1; i <= n; ++i)
       G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b) {
    G[a].pb(b), G[b].pb(a);
  void dfs(int u, int f, int d) {
    w[u] = 1, pa[u] = f, deep[u] = d++; for (int &i : G[u])
       if (i != f) +
         dfs(i, u, d), w[u] += w[i];
          if (w[mxson[u]] < w[i]) mxson[u] = i;
   \begin{array}{l} \mbox{void } \mbox{cut(int } \mbox{u, int } \mbox{link)} \ \{ \\ \mbox{data[pl[u] = ++t] = val[u], ulink[u] = link;} \end{array} 
    if (!mxson[u]) return;
    cut(mxson[u], link);
for (int i : G[u])
       if (i \neq pa[u] && i \neq mxson[u])
         cut(i, i);
  }
```

```
void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
        \label{eq:continuous_problem} \begin{array}{lll} \mbox{int} & \mbox{ta} = \mbox{ulink} \left[ \mbox{a} \right], & \mbox{tb} = \mbox{ulink} \left[ \mbox{b} \right], & \mbox{res} = \mbox{0}; \end{array}
        while (ta != tb) {
            if (deep
               `[ta] > deep[tb]) swap(ta, tb), swap(a, b);
'query(pl[tb], pl[b])
            tb = ulink[b = pa[tb]];
        if (pl[a] > pl[b]) swap(a, b);
        // query(pl[a], pl[b])
};
```

3.6

```
Centroid Decomposition* [5a24da]
\begin{array}{c} \textbf{struct} \;\; Cent\_Dec \; \{ \; // \;\; 1\text{-base} \\ vector < pll > G[N] \; ; \end{array}
  pll info[N]; // store info. of itself
pll upinfo[N]; // store info. of climbing up
int n, pa[N], layer[N], sz[N], done[N];
   ll \ dis [\underline{\phantom{a}} lg (N) + 1][N];
   void init(int _n) {
     n = _n, [ayer[0]] = -1;
      \begin{array}{l} \text{fill\_n} (pa + 1, n, 0), & \text{fill\_n} (done + 1, n, 0); \\ \text{for (int } i = 1; i <= n; ++i) & G[i].clear(); \end{array} 
   void add_edge(int a, int b, int w) {
     G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
   void get_cent(
     int u, int f, int &mx, int &c, int num) {
     int mxsz = 0;
     sz[u] = 1;
     for (pll e : G[u])

if (!done[e.X] && e.X != f) {
          if (mx > max(mxsz, num - sz[u]))
       mx = max(mxsz, num - sz[u]), c = u;
  if (!done[e.X] && e.X != f)
          dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
  int mx = 1e9, c = 0, lc;
     get_cent(u, f, mx, c, num);
     done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;

for (pll e : G[c])
        if (!done[e.X]) {
          if (sz[e.X] > sz[c])
          \begin{array}{l} lc = cut(e.X, c, num - sz[c]); \\ else \ lc = cut(e.X, c, sz[e.X]); \end{array}
          upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
        }
     return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
     for (int a = u, ly = layer[a]; a;
            a = pa[a],
                          --ly)
        info[a].X += dis[ly][u], ++info[a].Y;
        if (pa[a])
          upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
     }
  il query(int u) {
     11 \text{ rt} = 0;
     for (int a = u, ly = layer[a]; a;
        a = pa[a], --ly) {
rt += info[a].X + info[a].Y * dis[ly][u];
        if (pa[a])
             upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
     return rt;
};
```

LiChaoST* [4a4bee]

```
struct L {
```

if (x->f->isr()) rotate(x);

```
ll m, k, id;
                                                                               else if (x->dir() = x->f->dir())
  L() : id(-1) \{ \}
                                                                                  rotate(x->f), rotate(x);
  L(11 a, 11 b, 11 c) : m(a), k(b), id(c) {}
11 at(11 x) { return m * x + k; }
  L(ll a,
                                                                               else rotate(x), rotate(x);
                                                                          Splay* access(Splay *x) {
class LiChao { // maintain max
                                                                            Splay *q = nil;
private:
                                                                             for (; x != nil; x = x->f)
  int n; vector<L> nodes;
  void insert(int l, int r, int rt, L ln) {
                                                                               \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
     int m = (l + r) >> 1;
                                                                             return q;
     if (nodes[rt].id == -1)
       return nodes[rt] = ln, void();
                                                                          void root_path(Splay *x) { access(x), splay(x); }
     bool at Left = nodes[rt].at(1) < ln.at(1);
                                                                          void chroot(Splay *x){
     if (nodes[rt].at(m) < ln.at(m))
                                                                             root_path(x), x->give_tag(1);
       atLeft = 1, swap(nodes[rt], ln);
                                                                            x->push(), x->pull();
     if (r - l = 1) return;
                                                                          }
     if (atLeft) insert(l, m, rt \ll 1, ln);
                                                                          void split (Splay *x, Splay *y) {
     else insert (m, r, rt \ll 1 | 1, ln);
                                                                             chroot(x), root_path(y);
  11 query(int 1, int r, int rt, 11 x) {
  int m = (1 + r) >> 1; ll ret = -INF;
  if (nodes[rt].id != -1) ret = nodes[rt].at(x);
                                                                          void link(Splay *x, Splay *y) {
  root_path(x), chroot(y);
  x->setCh(y, 1);
     if (r - l = 1) return ret;
     if (x
                                                                          void cut(Splay *x, Splay *y) {
          < m) return max(ret, query(l, m, rt << 1, x));
                                                                             split(x, y);
                                                                             if (y->size != 5) return;
      return max(ret, query(m, r, rt << 1 | 1, x)); 
                                                                             y->push();
public:
                                                                            y->ch[0] = y->ch[0]->f = nil;
  Splay* get_root(Splay *x)
                                                                             for (root\_path(x); x->ch[0] != nil; x = x->ch[0])
  ll query(ll x) { return query(0, n, 1, x); }
                                                                               x \rightarrow push();
                                                                             splay(x);
3.8 Link cut tree* [a35b5d]
                                                                             return x;
struct Splay { // xor-sum
                                                                          bool conn(Splay *x, Splay *y) {
  static Splay nil;
Splay *ch[2], *f;
                                                                             return get\_root(x) == get\_root(y);
  int val, sum, rev, size;
                                                                          Splay* lca(Splay *x, Splay *y) {
  Splay (int
                                                                            access(x), root_path(y);
if (y->f == nil) return y;
        \underline{\text{val}} = 0) : val(\underline{\text{val}}), sum(\underline{\text{val}}), rev(0), size(1)
  \{ f = ch[0] = ch[1] = &nil; \}
                                                                             return y->f;
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
                                                                          void change(Splay *x, int val) {
  int dir()
                                                                            splay(x), x->val = val, x->pull();
  { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
                                                                          int query (Splay *x, Splay *y) {
     ch[d] = c;
                                                                            split(x, y);
     if (c != \&nil) c->f = this;
                                                                             return y->sum;
     pull();
                                                                          }
  void give_tag(int r) {
                                                                          3.9 KDTree [375ca2]
    if (r) swap(ch[0], ch[1]), rev = 1;
                                                                          namespace kdt {
  void push() {
                                                                          int root, lc [maxn], rc [maxn], xl [maxn], xr [maxn],
    if (ch[0] != &nil) ch[0]->give\_tag(rev);
if (ch[1] != &nil) ch[1]->give\_tag(rev);
                                                                             yl[maxn], yr[maxn];
                                                                          point p[maxn];
     rev = 0;
                                                                          int build(int l, int r, int dep = 0) {
                                                                             if (l = r) return -1;
                                                                             function < bool (const point &, const point &)> f =
  void pull() {
                                                                               [dep](const point &a, const point &b) {
  if (dep & 1) return a.x < b.x;</pre>
     // take care of the nil!
    size = ch[0]->size + ch[1]->size + 1;

sum = ch[0]->sum ^ ch[1]->sum ^ val;

if (ch[0]!= &nil) ch[0]->f = this;

if (ch[1]!= &nil) ch[1]->f = this;
                                                                                  else return a.y < b.y;
                                                                             int m = (1 + r) >> 1;
                                                                            } Splay::nil;
Splay *nil = &Splay::nil;
                                                                             yl[m] = yr[m] = p[m].y;
void rotate (Splay
                                                                             lc[m] = build(1, m, dep + 1);
                       *x) {
  Splay *p = x - f;
                                                                             if (~lc[m]) {
                                                                                \begin{array}{l} xl\left[m\right] = \min(xl\left[m\right], \ xl\left[lc\left[m\right]\right]); \\ xr\left[m\right] = \max(xr\left[m\right], \ xr\left[lc\left[m\right]\right]); \\ \end{array} 
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
                                                                               yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
  else x->f = p->f;
  p-\operatorname{setCh}(x-\operatorname{sch}[!d], d);
  x->setCh(p, !d);
  p->pull(), x->pull();
                                                                             rc[m] = build(m + 1, r, dep + 1);
                                                                             if (\sim rc [m]) {
void splay(Splay *x) {
                                                                               xl[m] = min(xl[m], xl[rc[m]]);
                                                                                \begin{array}{ll} xr\left[m\right] &= \max(xr\left[m\right], \ xr\left[rc\left[m\right]\right]); \\ yl\left[m\right] &= \min(yl\left[m\right], \ yl\left[rc\left[m\right]\right]); \end{array} 
  vector < Splay*> splayVec;
  for (Splay *q = x;; q = q->f) {
     splayVec.pb(q);
                                                                               yr[m] = max(yr[m], yr[rc[m]]);
     if (q->isr()) break;
                                                                             return m;
  reverse (ALL(splayVec));
  for (auto it : splayVec) it->push();
                                                                          bool bound(const point &q, int o, long long d) {
  while (!x->isr()) {
                                                                             double ds = sqrt(d + 1.0);
```

if (q.x < xl[o] - ds || q.x > xr[o] + ds ||

```
q.y < yl[o] - ds | | q.y > yr[o] + ds)
    return false;
                                                                           node *\&t = k < o->data ? o->l : o->r;
  return true;
                                                                           return erase(t, k) ? o->up(), 1 : 0;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
  (a.y - b.y) * 111 * (a.y - b.y);
                                                                        void insert(node *&o, int k) {
                                                                           node *a, *b;
                                                                           split(o, a, b, k),
                                                                             o = merge(a, merge(new node(k), b));
void dfs(
                                                                        void interval(node *&o, int l, int r) {
  const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
long long cd = dist(p[o], q);
                                                                           node *a, *b, *c;
split2(o, a, b, l - 1), split2(b, b, c, r);
  if (cd != 0) d = min(d, cd);
if ((dep \& 1) \&\& q.x < p[o].x
                                                                           // operate
                                                                           o = merge(a, merge(b, c));
    4 Flow/Matching
  } else {
                                                                        4.1 Dinic [98fb3a]
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                        {\color{red} \textbf{struct}} \hspace{0.1cm} \textbf{MaxFlow} \hspace{0.1cm} \{ \hspace{0.1cm} // \hspace{0.1cm} \textbf{0-base}
  }
                                                                           struct edge {
                                                                             int to, cap, flow, rev;
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
                                                                           vector<edge> G[MAXN];
  root = build(0, v.size());
                                                                           \label{eq:cur_maxn} \begin{array}{lll} \textbf{int} & s \;, \; t \;, \; dis \left[ \text{MAXN} \right] \;, \; cur \left[ \text{MAXN} \right] \;, \; n \,; \end{array}
                                                                           int dfs(int u, int cap) {
  if (u == t || !cap) return cap;
long long nearest (const point &q) {
  long long res = 1e18;
                                                                              for (int &i = cur[u]; i < (int)G[u]. size(); ++i) {
  dfs(q, res, root);
                                                                                edge \&e = G[u][i]
  return res;
                                                                                if (dis[e.to] = dis[u] + 1 & e.flow != e.cap)
                                                                                   int df = dfs(e.to, min(e.cap - e.flow, cap));
} // namespace kdt
                                                                                   if (df) {
3.10 Treap [5ab1a1]
                                                                                     e.flow += df;
G[e.to][e.rev].flow -= df;
struct node {
                                                                                     return df;
  int data, sz;
node *l, *r;
                                                                                   }
                                                                                }
  node(int k) : data(k), sz(1), l(0), r(0) \{ \}
  void up() {
                                                                              dis[u] = -1;
    sz = 1;
                                                                              return 0:
     \begin{array}{lll} \textbf{if} & (\ l\ ) & \text{sz} \ +\!\!= \ l\!-\!\!>\! \text{sz} \ ; \\ \end{array} 
    if (r) sz += r->sz;
                                                                           bool bfs() {
                                                                              fill_n(dis, n, -1);
  void down() {}
                                                                             queue<int> q;
q.push(s), dis[s] = 0;
};
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
                                                                              while (!q.empty()) {
                                                                                int tmp = q. front();
  if (!a | | !b) return a ? a : b;
                                                                                q.pop();
  if (rand()\%(sz(a) + sz(b)) < sz(a))
                                                                                for (auto &u : G[tmp])
    if (!~dis[u.to] & u.flow != u.cap) {
                                                                                     q.push(u.to);
  return b->down(), b->l = merge(a, b->l), b->up(), b;
                                                                                      dis[u.to] = dis[tmp] + 1;
void split (node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
                                                                              return dis[t] != -1;
  o > down();
  if (o->data \le k)
                                                                           int maxflow(int _s, int _t) {
    a = o, split(o->r, a->r, b, k), a->up();
                                                                              s = \_s, t = \_t;
  else b = o, split (o->1, a, b->1, k), b->up();
                                                                              int flow = 0, df;
                                                                              while (bfs()) {
void split2 (node *o, node *&a, node *&b, int k) {
                                                                                fill_n(cur, n, 0);
  if (sz(o) \le k) return a = o, b = 0, void();
                                                                                while ((df = dfs(s, INF))) flow += df;
  o->down();
  if (sz(o->1) + 1 \le k)
                                                                             return flow;
   \begin{array}{l} a = o, \; split2\,(o->r\,,\; a->r\,,\; b,\; k\; -\; sz\,(o->l)\; -\; 1)\,; \\ else\; b = o, \; split2\,(o->l\,,\; a,\; b->l\,,\; k)\,; \end{array} 
                                                                           void init(int _n) {
  o > up();
                                                                              for (int i = 0; i < n; ++i) G[i].clear();
node *kth(node *o, int k) {
  if (k \le sz(o->1)) return kth(o->1, k);
if (k = sz(o->1) + 1) return o;
                                                                           void reset() {
                                                                              for (int i = 0; i < n; ++i)
  return kth(o->r, k - sz(o->l) - 1);
                                                                                for (auto \& j : G[i]) j.flow = 0;
int Rank(node *o, int key) {
                                                                           void add_edge(int u, int v, int cap) {
   G[u].pb(edge{v, cap, 0, (int)G[v].size()});
  if (!o) return 0;
if (o->data < key)</pre>
                                                                             G[v].pb(edge\{u, 0, 0, (int)G[u].size() - 1\});
    return sz(o->l) + 1 + Rank(o->r, key);
  else return Rank(o->l, key);
bool erase(node *&o, int k) {
                                                                              Bipartite Matching* [784535]
                                                                        4.2
  if (!o) return 0;
  if (o->data == k) {
                                                                        struct Bipartite_Matching { // 0-base
    node *t = o;
                                                                           int mp[N], mq[N], dis[N+1], cur[N], l, r;
    o->down(), o = merge(o->l, o->r);
                                                                           vector < int > G[N + 1];
                                                                           bool dfs(int u) {
    delete t;
    return 1;
                                                                              for (int &i = cur[u]; i < SZ(G[u]); ++i) {
```

```
National Yang Ming Chiao Tung University FubukiMyWife
       int e = G[u][i];
       return mp[mq[e] = u] = e, 1;
    \mathbf{return} \ \mathbf{dis} [\mathbf{u}] = -1, \ 0;
  bool bfs() {
    queue<int> q;
    fill_n (dis, l + 1, -1);
for (int i = 0; i < l; ++i)
       if (! ~mp[i])
         q.push(i), dis[i] = 0;
     while (!q.empty()) {
       int \dot{\mathbf{u}} = \mathbf{q}. front();
       q.pop();
       for (int e : G[u])
         if (!~dis[mq[e]])
           q.\,push\,(mq[\,e\,]\,)\ ,\ dis\,[mq[\,e\,]\,]\ =\ dis\,[\,u\,]\ +\ 1;
    return dis[1] != -1;
  int matching() {
    int res = 0;
    fill_n(mp, l, -1), fill_n(mq, r, l);
    while (bfs()) {
       fill_n (cur, 1, 0);
       for (int i = 0; i < l; ++i)
         res += (! \sim mp[i] \&\& dfs(i));
    return res; // (i, mp[i] != -1)
  void add_edge(int s, int t) { G[s].pb(t); }
  void init(int _l, int _r) {
    l = _l, r = _r;
for (int i = 0; i <= 1; ++i)
      G[i].clear();
4.3 Kuhn Munkres* [4b3863]
struct KM {
             // 0-base, maximum matching
  11 w[N] [N], h1 [N], hr [N], s1k [N]; int f1 [N], fr [N], pre [N], qu [N], ql, qr, n; bool v1 [N], vr [N];
  void init(int _n) {
    n\,=\,\underline{\phantom{a}}n;
    for (int i = 0; i < n; +++i)
       fill_n (w[i], n, -INF);
```

```
void add_edge(int a, int b, ll wei) {
  w[a][b] = wei;
bool Check(int x) {
  if (vl[x] = 1, \sim fl[x])
  return \operatorname{vr}[\operatorname{qu}[\operatorname{qr}++] = \operatorname{fl}[x]] = 1;
while (\sim x) \operatorname{swap}(x, \operatorname{fr}[\operatorname{fl}[x] = \operatorname{pre}[x]]);
  return 0;
void bfs(int s) {
  fill_n(slk
       , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
  ql = qr = 0, qu[qr++] = s, vr[s] = 1;
  for (ll d;;) {
     while (q\hat{l} < qr)
for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
    d = INF;
     for (int x = 0; x < n; ++x)
       if (!vl[x] \& d > slk[x])' d = slk[x];
     for (int x = 0; x < n; ++x) {
       if(vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (\operatorname{vr}[x]) \operatorname{hr}[x] = d;
     for (int x = 0; x < n; ++x)
        if (!vl[x] && !slk[x] && !Check(x)) return;
  }
ll solve() {
  fill_n (fl
       , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
```

```
for (int i = 0; i < n; ++i)
    hl[i] = *max_element(w[i], w[i] + n);
    for (int i = 0; i < n; ++i) bfs(i);
    ll res = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];
    return res;
}
};</pre>
```

4.4 MincostMaxflow* [1c78db]

```
struct MinCostMaxFlow { // 0-base
   struct Edge {
      ll from, to, cap, flow, cost, rev;
   } *past[N];
    vector<Edge> G[N];
   int inq[N], n, s, t;
ll dis[N], up[N], pot[N];
bool BellmanFord() {
      fill_n(dis, n, INF), fill_n(inq, n, 0);
      queue < int > q;
      auto relax = [&](int u, ll d, ll cap, Edge *e) {
    if (cap > 0 && dis[u] > d) {
            dis[u] = d, up[u] = cap, past[u] = e;
            if (!inq[u]) inq[u] = 1, q.push(u);
         }
      };
      relax(s, 0, INF, 0);
while (!q.empty()) {
        int u = q.front();
         q.pop(), inq[u] = 0;
for (auto &e : G[u]) {
           11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
                 (e.to, d2, min(up[u], e.cap - e.flow), &e);
        }
      }
      return dis[t] != INF;
   void solve (int
      , int _t, ll &flow, ll &cost, bool neg = true) { s = \_s, t = \_t, flow = 0, cost = 0;
      if (neg) BellmanFord(), copy_n(dis, n, pot);
      for (; BellmanFord(); copy_n(dis, n, pot)) {
         for (int
         i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
         for (int i = t; past[i]; i = past[i]->from) {
           auto &e = *past[i];
            e.flow += up[t], G[e.to][e.rev].flow -= up[t];
         }
      }
   void init(int _n) {
      n = n, fill_n(pot, n, 0);
for (int i = 0; i < n; ++i) G[i].clear();
    void add_edge(ll a, ll b, ll cap, ll cost)
      \begin{array}{l} G[a].p\overline{b}(Edge\{a,\ b,\ cap,\ 0,\ cost\ ,\ SZ(G[b])\ \})\,;\\ G[b].pb(Edge\{b,\ a,\ 0,\ 0,\ -cost\ ,\ SZ(G[a])\ -\ 1\})\,; \end{array}
};
```

4.5 Maximum Simple Graph Matching* [0fe1c3]

```
struct Matching { // 0-base
  queue<int> q; int n;
   vector<int> fa, s, vis, pre, match;
   vector<vector<int>>> G;
  int Find(int u)
   \{ return u = fa[u] ? u : fa[u] = Find(fa[u]); \}
  int LCA(int x, int y) {
     \begin{array}{lll} \textbf{static} & \textbf{int} & tk = 0; & tk++; & x = Find(x); & y = Find(y); \end{array}
     for (;; swap(x, y)) if (x != n) {
       if (vis[x] = tk) return x;
       vis[x] = tk;
       x = Find(pre[match[x]]);
    }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != 1; x = pre[y]) {
       pre[x] = y, y = match[x];
       if (s[y] = 1) q.push(y), s[y] = 0;
for (int z: \{x, y\}) if (fa[z] = z) fa[z] = 1;
     }
  bool Bfs(int r) {
```

```
\begin{array}{ll} iota\left(ALL(fa)\,,\;\;0\right);\;\; fill\left(ALL(s)\,,\;\;\text{-1}\right);\\ q=queue<int>();\;\;q.push(r)\,;\;\;s\left[r\right]\,=\,0; \end{array}
   for (; !q.empty(); q.pop()) {
            (int x = q. front(); int u : G[x])
          if(s[u] = -1) {
             if (pre[u] = x, s[u] = 1, match[u] == n) {
for (int a = u, b = x, last;
                      b != n; a = last, b = pre[a]
                         match[b], match[b] = a, match[a] = b;
                return true;
             q.push(match[u]); s[match[u]] = 0;
         } else if (!s[u] \&\& Find(u) != Find(x)) {
             int l = LCA(u, x);
             Blossom\left(x,\ u,\ l\right);\ Blossom\left(u,\ x,\ l\right);
   return false;
Matching(\, \underline{i}\, \underline{n} t \,\, \underline{\hspace{1em}} n) \,\, : \,\, n(\underline{\hspace{1em}} n) \,\, , \,\, fa\,(n\, +\, 1) \,, \,\, s\,(n\, +\, 1) \,,
\begin{array}{c} (n+1)\,,\;pre\,(n+1,\,n)\,,\;match\,(n+1,\,n)\,,\;G(n)\ \{\}\\ void\ add\_edge\,(int\ u,\ int\ v) \end{array}
\{G[u].pb(v), G[v].pb(u); \}
int solve() {
   int ans = 0;
   for (int x = 0; x < n; ++x)
        if (match[x] == n) ans += Bfs(x); 
   return ans:
} // match[x] == n means not matched
```

4.6 Maximum Weight Matching* [1ec446]

```
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
   struct edge { int u, v, w; }; int n, nx;
   vector<int> lab; vector<vector<edge>>> g;
   \label{eq:vector} {\tt vector}{<} {\tt int}{\gt} \ {\tt slk} \ , \ \ {\tt match} \ , \ \ {\tt st} \ , \ \ {\tt pa} \ , \ \ {\tt S} \ , \ \ {\tt vis} \ ;
   vector(\operatorname{int} > \operatorname{sir}, \operatorname{match}, \operatorname{sir}, \operatorname{pa}, \operatorname{b}, \operatorname{vis}, vector(\operatorname{vector} < \operatorname{int} > \operatorname{flo}, \operatorname{flo}, \operatorname{flo}, \operatorname{from}; \operatorname{queue} < \operatorname{int} > \operatorname{q}; WeightGraph(\operatorname{int} \operatorname{n}_{-}): \operatorname{n}(\operatorname{n}_{-}), \operatorname{nx}(\operatorname{n} * 2), \operatorname{lab}(\operatorname{nx} + 1), \operatorname{g}(\operatorname{nx} + 1, \operatorname{vector} < \operatorname{edge} > (\operatorname{nx} + 1)), \operatorname{slk}(\operatorname{nx} + 1),
       flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
      match = st = pa = S = vis = slk;
      REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
   int E(edge e)
   \{ return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; \}
   void set_slk(int x) {
       slk[x] = 0;
      REP(u, 1, n)
          if (g[u][x].w > 0 \& st[u] != x \& S[st[u]] == 0)
             update_slk(u, x, slk[x]);
   void q_push(int x) {
      if (x \le n) q.push(x);
      else for (int y : flo [x]) q_push(y);
   void set_st(int x, int b) {
      st[x] = b;
      if (x > n) for (int y : flo[x]) set_st(y, b);
   vector<int> split_flo(auto &f, int xr) {
      auto it = find(ALL(f), xr);
if (auto pr = it - f.begin(); pr % 2 == 1)
          reverse(1 + ALL(f)), it = f.end() - pr;
      auto res = vector(f.begin(), it);
      return f.erase(f.begin(), it), res;
   void set_match(int u, int v) {
      match[u] = g[u][v].v;
       if (u \le n) return;
      int xr = flo_from[u][g[u][v].u];
      set_match(xr, v); f.insert(f.end(), ALL(z));
   void augment(int u, int v) {
      for (;;) {
          \begin{array}{ll} \textbf{int} & \textbf{xnv} = \, \textbf{st} \, [\, \textbf{match} \, [\, \textbf{u} \, ] \, ] \, ; \  \, \textbf{set\_match} \, (\, \textbf{u} \, , \  \, \textbf{v} \, ) \, ; \end{array}
          if (!xnv) return;
          set_{match}(v = xnv, u = st[pa[xnv]]);
   int lca(int u, int v) {
```

```
\begin{array}{lll} \textbf{static} & \textbf{int} & \textbf{t} = \ \textbf{0}; \ \textbf{++}\textbf{t}; \end{array}
   \quad \text{for } (++t \; ; \; u \; | \; | \; | \; v \; ; \; swap(u, \; v)) \; \; if \; \; (u) \; \; \{
     if (vis[u] == t) return u;
vis[u] = t, u = st[match[u]];
     if (u) u = st[pa[u]];
   return 0;
void add_blossom(int u, int o, int v) {
   int b = find(n + 1 + ALL(st), 0) - begin(st);
   lab[b] = 0, \dot{S}[b] = 0, match[b] = match[o];
   vector < int > f = \{o\};
   for (int t : {u, v}) {
      reverse(1 + ALL(f));
     for (int x = t, y; x != o; x = st[pa[y]])
f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set_st(b, b);

REP(x, 1, nx) g[b][x].w = g[x][b].w = 0;

fill (ALL(flo_from[b]), 0);
   for (int xs : flo[b]) {
     REP(x, 1, nx)
         if (g[b][x].w = 0 \mid | E(g[xs][x]) < E(g[b][x])
           g[b][x] = g[xs][x], g[x][b] = g[x][xs];
     REP(x, 1, n)
         if (flo_from [xs][x]) flo_from [b][x] = xs;
   set_slk(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u],
   for (int x : split_flo(flo[b], xr)) {
     if (xs = -1) { xs = x; continue; } pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
      slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
   for (int x : flo[b])
      if (x = xr) S[x] = 1, pa[x] = pa[b];
      else S[x] = -1, set\_slk(x);
   st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] = -1) {
    int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
      slk[v] = slk[nu] = S[nu] = 0; q_push(nu);
   else if (S[v] = 0)
      if (int o = lca(u, v)) add_blossom(u, o, v);
      else return augment(u, v), augment(v, u), true;
   return false;
bool matching() {
   fill (ALL(S), -1), fill (ALL(slk), 0);
q = queue<int>();
  \begin{array}{l} \text{REP}(x, 1, nx) \ \ \text{if} \ \ (\text{st}[x] = x \ \&\& \ ! \text{match}[x]) \\ \text{pa}[x] = S[x] = 0, \ \text{q_push}(x); \end{array}
   if (q.empty()) return false;
   for (;;)
      while (SZ(q)) {
         int u = q.front(); q.pop();
if (S[st[u]] == 1) continue;
        REP(v, 1, n)
            if (g[u][v].w > 0 & st[u] != st[v]) {
if (E(g[u][v]) != 0)
                  update_slk(u, st[v], slk[st[v]]);
                       (on_found_edge(g[u][v])) return true;
            }
      int d = INF;
     REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
     REP(x,\ 1,\ nx)
         if (int
                s = slk[x]; st[x] == x & s & s & s[x] <= 0
            d = min(d, E(g[s][x]) / (S[x] + 2));
     REP(u, 1, n)
         \begin{array}{l} \mbox{if } (S[st[u]] == 1) \mbox{ lab}[u] +\!\!\!= d; \\ \mbox{else if } (S[st[u]] =\!\!\!= 0) \mbox{ } \{ \end{array}
            if \ (lab\,[u] \mathrel{<=} d) \ return \ false\,;
            lab[u] -= d;
      \begin{array}{lll} REP(b, \ n+1, \ nx) & if \ (st[b] \Longrightarrow b \ \&\& \ S[b] >= 0) \\ lab[b] & += d \ * \ (2 \ -4 \ * \ S[b]) \ ; \end{array} 
     REP(x, 1, nx)
```

```
National Yang Ming Chiao Tung University FubukiMyWife
                                     if (on_found_edge(g[s][x])) return true;
                           REP(b, n + 1, nx)

if (st[b] = b && S[b] = 1 && lab[b] = 0)
                                             expand_blossom(b);
                    return false;
          pair<ll, int> solve() {
                    fill (ALL(match), 0);
                  REP(u, 0, n) st[u] = u, flo[u].clear();
                    int w_max = 0
                 \begin{array}{lll} & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\
                           w_max = max(w_max, g[u][v].w);
                    fill (ALL(lab), w_max);
                    int n_matches = 0; ll tot_weight = 0;
                  \begin{array}{c} \text{while } (\text{matching}\,()) + \text{h\_matches}\,; \\ \text{REP}(u,\ 1,\ n) \ \ \text{if} \ \ (\text{match}\,[u] \ \&\& \ \text{match}\,[u] \ < \ u) \\ \end{array} 
                            tot\_weight += g[u][match[u]].w;
                    return make_pair(tot_weight, n_matches);
          void add_edge(int u, int v, int w)
          \{g[u][v].w = g[v][u].w = w; \}
                          SW-mincut [6621f5]
 4.7
```

```
\begin{array}{ll} \textbf{int} & \text{vst} \left[ \textbf{MXN} \right] \,, & \text{edge} \left[ \textbf{MXN} \right] \left[ \textbf{MXN} \right] \,, & \text{wei} \left[ \textbf{MXN} \right] ; \end{array}
void init(int n)
  REP fill_n (edge[i], n, 0);
void addEdge(int u, int v, int w){
   edge[u][v] += w; edge[v][u] += w;
      search(int &s, int &t, int n){
int
   fill_n(vst, n, 0), fill_n(wei, n, 0);
   s = t = -1;
   int mx, cur;
   for (int j = 0; j < n; ++j) {
mx = -1, cur = 0;
      \label{eq:representation} \text{REP if } (\, wei \, [\, i \, ] \, > \, mx) \  \, cur \, = \, i \, , \, \, mx \, = \, wei \, [\, i \, ] \, ;
      vst[cur] = 1, wei[cur] = -1;
      s = t; t = cur;
      REP if (!vst[i]) wei[i] += edge[cur][i];
   return mx;
int solve(int n) {
   int res = INF;
```

 $\begin{array}{l} res = min(res, search(x, y, n)); \\ REP \ edge[i][x] = (edge[x][i] += edge[y][i]); \\ \end{array}$

BoundedFlow*(Dinic*) [4a793f]

edge[y][i] = edge[n - 1][i]; edge[i][y] = edge[i][n - 1];

for (int x, y; n > 1; n--){

return res;

} } sw;

```
struct BoundedFlow { // 0-base
   struct edge {
     \quad \text{int to, cap, flow, rev;} \\
   vector<edge> G[N]:
   int n, s, t, dis[N], cur[N], cnt[N];
   void init(int _n) {
     \begin{array}{lll} n = \underline{\ } n; \\ \text{for (int } i = 0; \ i < n + 2; +\!\!\!+\!\!\! i) \end{array}
        G[i].clear(), cnt[i] = 0;
   void add_edge(int u, int v, int lcap, int rcap) {
     cnt[u] = lcap, cnt[v] += lcap
     \begin{array}{l} G[u].pb(edge\{v,\ rcap\ ,\ lcap\ ,\ SZ(G[v])\ \})\,;\\ G[v].pb(edge\{u,\ 0\ ,\ 0\ ,\ SZ(G[u])\ -\ 1\})\,; \end{array}
   void add_edge(int u, int v, int cap) {
     G[u].pb(edge\{v, cap, 0, SZ(G[v])\});
```

```
G[v].pb(edge\{u, 0, 0, SZ(G[u]) - 1\});
  edge \&e = G[u][i];
       if (dis[e.to] = dis[u] + 1 & e.cap!= e.flow) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df, G[e.to][e.rev].flow -= df;
           return df;
         }
       }
     dis[u] = -1;
    return 0;
  bool bfs() {
     fill_n(dis, n + 3, -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (edge &e : G[u])
         if (!~dis[e.to] & e.flow != e.cap)
            q.push(e.to), dis[e.to] = dis[u] + 1;
     return dis[t] != -1;
  int maxflow(int _s, int _t) {
     s = 
          _{s}, t = _{t};
     int flow = 0, df;
     while (bfs()) {
       fill_n(cur, n + 3, 0);
       while ((df = dfs(s, INF))) flow += df;
     return flow;
  bool solve() {
     int sum = 0;
     for (int i = 0; i < n; ++i)
       if (cnt[i] > 0)
       \begin{array}{lll} & \text{add\_edge}(n+1,\ i,\ cnt[\,i\,])\ ,\ sum\ +=\ cnt[\,i\,]\,;\\ & \text{else if } (cnt[\,i\,]\ <\ 0)\ add\_edge(\,i\,,\ n+2\,,\ -cnt[\,i\,])\,; \end{array}
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)
       if (cnt[i] > 0)
       G[n + 1].pop_back(); G[i].pop_back(); else if (cnt[i] < 0)
         G[i].pop\_back(), G[n + 2].pop\_back();
     return sum != -1;
  int solve(int _s, int _t) {
  add_edge(_t, _s, INF);
    if (!solve()) return -1; // invalid flow int x = G[\_t]. back().flow;
     return G[_t].pop_back(), G[_s].pop_back(), x;
};
4.9 Gomory Hu tree* [11be99]
```

```
MaxFlow Dinic;
 int g [MAXN];
 void GomoryHu(int n) { // 0-base
     fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
         Dinic.reset();
        \begin{array}{l} {\rm add\_edge(i\,,\,\,g[\,i\,]\,,\,\,Dinic\,.\,maxflow(i\,,\,\,g[\,i\,])\,)\,;} \\ {\rm for\,\,(int\,\,j\,=\,i\,+\,1;\,\,j\,<=\,n;\,\,+\!\!+\!\!j\,)} \end{array}
            if (g[j] == g[i] &  \sim Dinic.dis[j])
                g[j] = i;
}
```

4.10 Minimum Cost Circulation* [ba97cf]

```
struct MinCostCirculation { // 0-base
  struct Edge {
     ll\ from\,,\ to\,,\ cap\,,\ fcap\,,\ flow\,,\ cost\,,\ rev\,;
  } *past[N];
  vector < Edge > G[N];
  11 dis[N], inq[N], n;
void BellmanFord(int s) {
     fill_n(dis, n, INF), fill_n(inq, n, 0);
     queue<int> q;
```

```
auto relax = [&](int u, ll d, Edge *e) {
  if (dis[u] > d) {
    dis[u] = d, past[u] = e;
}
           if (!inq[u]) inq[u] = 1, q.push(u);
       }
     };
     \operatorname{relax}\left(s\,,\;\;0\,,\;\;0\right);
     while (!q.empty()) {
        int u = q.front();
        q.pop(), inq[u] = 0;
        for (auto &e: G[u])
if (e.cap > e.flow)
             relax(e.to, dis[u] + e.cost, \&e);
   }
   void try_edge(Edge &cur) {
     if (cur.cap > cur.flow) return ++cur.cap, void();
     BellmanFord(cur.to);
     if (dis[cur.from] + cur.cost < 0) {
    ++cur.flow, --G[cur.to][cur.rev].flow;</pre>
        for (int
              i = cur.from; past[i]; i = past[i]->from) {
          auto &e = *past[i];
++e.flow, --G[e.to][e.rev].flow;
     ++cur.cap;
   void solve(int mxlg) {
     for (int b = mxlg; b >= 0; --b) {
  for (int i = 0; i < n; ++i)
  for (auto &e : G[i])
             e.cap *= 2, e.flow *= 2;
        for (int i = 0; i < n; ++i)
           for (auto &e : G[i])
             if (e.fcap >> b & 1)
               try_edge(e);
     }
   void init (int _n) { n = _n;
     for (int i = 0; i < n; ++i) G[i].clear();
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].pb(Edge
           \{a, b, 0, cap, 0, cost, SZ(G[b]) + (a = b)\}\);
     G[b].pb(Edge\{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1\});
} mcmf; // O(VE * ElogC)
```

4.11 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct supersource S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - − To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - − To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite $\operatorname{graph}(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
- 1. Consruct super source S and $\sinh T$
- 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) = (0, d(v))
- 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - $1. \ \ {\bf Binary search \, on \, answer, suppose \, we're \, checking \, answer} \, T$
- 2. Construct a max flow model, let K be the sum of all weights

- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
- 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
- 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$
 $p_{u} \ge 0$

5 String

5.1 KMP [5a0728]

```
 \begin{array}{l} & \text{int } F[\text{MAXN}]; \\ & \text{vector} < \text{int} > \text{match} (\text{string } A, \text{ string } B) \  \, \{ \\ & \text{vector} < \text{int} > \text{ans}; \\ & F[\,0\,] = -1, \  \, F[\,1\,] = \, 0; \\ & \text{for } (\text{int } i = 1, \ j = 0; \  \, i < SZ(B); \  \, F[++i\,] = ++j) \  \, \{ \\ & \text{if } (B[\,i\,] = B[\,j\,]) \  \, F[\,i\,] = F[\,j\,]; \  \, // \  \, \text{optimize} \\ & \text{while } (\,j\,! = -1 \, \&\& \, B[\,i\,] \, ! = \, B[\,j\,]) \  \, j = \, F[\,j\,]; \\ & \} \\ & \text{for } (\text{int } i = 0, \ j = 0; \  \, i < SZ(A); \  \, +\!+i) \  \, \{ \\ & \text{while } (\,j\,! = -1 \, \&\& \, A[\,i\,] \, ! = \, B[\,j\,]) \  \, j = \, F[\,j\,]; \\ & \text{if } (++j = \, SZ(B)) \  \, \text{ans.pb}(\,i\,+1\,-\,j\,), \  \, j = \, F[\,j\,]; \\ & \} \\ & \text{return } \, \text{ans}; \\ & \} \\ \end{array}
```

5.2 Z-value* [b47c17]

6.3 Manacher* [1ad8ef]

5.4 SAIS* [6f26bc]

```
auto sais(const auto &s) { const int n = SZ(s), z = ranges::max(s) + 1; if (n = 1) return vector\{0\}; vector<int>c(z); for (int x : s) ++c[x]; partial_sum(ALL(c), begin(c)); vector<int>sa(n); auto I = views::iota(0, n); vector<bool>t(n, true);
```

```
for (int i = n - 2; i >= 0; --i)
     t[i] = (
  s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
auto is_lms = views::filter([&t](int x) {
     return x && t[x] && !t[x - 1];
  });
  auto induce = [&] {
     for (auto x = c; int y : sa)

if (y--) if (!t[y]) sa[x[s[y]-1]++] = y;
     for (auto x = c; int y : sa | views::reverse)
        if (y--) if (t[y]) sa[--x[s[y]]] = y;
  vector<int> lms, q(n); lms.reserve(n);
for (auto x = c; int i : I | is_lms)
    q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
  induce(); vector<int> ns(SZ(lms));
  for (int j = -1, nz = 0; int i : sa \mid is\_lms) {
     if (j >= 0) {
        int len = min({n - i, n - j, lms[q[i] + 1] - i});
        ns[q[i]] = nz += lexicographical_compare(
             \begin{array}{lll} \text{begin}(s) + j, & \text{begin}(s) + j + \text{len}, \\ \text{begin}(s) + i, & \text{begin}(s) + i + \text{len}); \end{array}
     j = i;
  fill(ALL(sa), 0); auto nsa = sais(ns);
for (auto x = c; int y : nsa | views::reverse)
     y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
// sa[i]: sa[i]-th suffix
       is the i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i-1].
struct Suffix {
  int n; vector<int> sa, hi, ra;
  Suffix
        (const auto \&\_s, int \_n) : n(\_n), hi(n), ra(n) 
     vector<int> s(n + 1); // s[n] = 0; copy_n(_s, n, begin(s)); // _s shouldn't contain 0 sa = sais(s); sa.erase(sa.begin()); for (int i = 0; i < n; ++i) ra[sa[i]] = i; for (int i = 0, h = 0; i < n; ++i) {
        if (!ra[i]) { h = 0; continue; }
        for (int j = sa[ra[i] - 1]; max
              (i, j) + h < n & s[i + h] = s[j + h];) + h;
        hi[ra[i]] = h ? h-- : 0;
};
        Aho-Corasick Automatan* [794a77]
struct AC_Automatan {
```

```
int nx[len][sigma],
                         fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
int newnode() {
  fill_n(nx[top], sigma, -1);
  return top++;
void init() { top = 1, newnode(); }
int input(string &s) {
  int X = 1;
  \begin{array}{lll} & for \;\; (char \;\; c \;\; : \; s) \;\; \{ \\ & if \;\; (!\!\sim\! nx [X] [\; c \;\; - \;\; 'A']) \;\; nx [X] [\; c \;\; - \;\; 'A'] \; = \; newnode() \; ; \end{array}
     X = nx[X][c - A'];
  return X; // return the end node of string
void make_fl() {
  queue<int> q;
  q.push(1), fl[1] = 0;
for (int t = 0; !q.empty(); ) {
     int R = q.front();
     q.pop(), ord[t++] = R;
for (int i = 0; i < sigma; ++i)
       if (~nx[R][i])
          else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
  }
void solve() {
  for (int i = top - 2; i > 0; --i)
     cnt[fl[ord[i]]] += cnt[ord[i]];
```

5.6 Smallest Rotation [4f469f]

} ac;

```
\begin{array}{l} string \ mcp(string \ s) \ \{\\ int \ n = SZ(s) \,, \ i = 0 \,, \ j = 1 \,;\\ s += s \,;\\ while \ (i < n \,\&\& \, j < n) \, \{\\ int \ k = 0 \,;\\ while \ (k < n \,\&\& \, s[\, i + k] == s[\, j + k]) \ ++k \,;\\ if \ (s[\, i + k] <= s[\, j + k]) \ j \,+= k \,+ 1 \,;\\ else \ i \,+= k \,+ 1 \,;\\ if \ (i == j) \,++j \,;\\ \}\\ int \ ans = i < n \,\,? \,\, i \,: \,\, j \,;\\ return \ s.substr(ans \,, \, n) \,;\\ \} \end{array}
```

5.7 De Bruijn sequence* [a09470]

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N
  void dfs(int *out, int t, int p, int &ptr) {
    if (ptr >= L) return;
    if (t > N) {
        if (N % p) return;
        for (int i = 1; i <= p && ptr < L; ++i)
            out[ptr++] = buf[i];
    } else {
        buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
        for (int j = buf[t - p] + 1; j < C; ++j)
            buf[t] = j, dfs(out, t + 1, t, ptr);
        }
    }
}
void solve(int _c, int _n, int _k, int *out) {
    int p = 0;
    C = _c, N = _n, K = _k, L = N + K - 1;
        dfs(out, 1, 1, p);
        if (p < L) fill(out + p, out + L, 0);
    }
} dbs;</pre>
```

5.8 Extended SAM* [64c3b7]

```
struct exSAM {
  int len[N * 2], link[N * 2]; // maxlength, suflink
  int next[N * 2][CNUM], tot; // [0, tot), root = 0
  int lenSorted[N * 2]; // topo. order
   int cnt[N * 2]; // occurence
   int newnode() {
      fill_n (next[tot], CNUM, 0);
      len[tot] = cnt[tot] = link[tot] = 0;
      return tot++;
   void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
      int cur = next[last][c];
      len[cur] = len[last] + 1;
     int p = link[last];
while (p != -1 && !next[p][c])
   next[p][c] = cur, p = link[p];
      if (p = -1) return link[cur] = 0, cur;
      int q = next[p][c];
      if (len
            [p] + 1 \Longrightarrow len[q]) return link[cur] = q, cur;
      int clone = newnode():
      for (int i = 0; i < CNUM; ++i)
        next[
              clone [[i] = len [next[q][i]] ? next[q][i] : 0;
     \begin{array}{l} len[clone] = len[p] + 1; \\ while & (p != -1 \&\& next[p][c] == q) \end{array}
      next[p][c] = clone, p = link[p];
link[link[cur] = clone] = link[q];
      link[q] = clone;
      return cur;
   void insert(const string &s) {
      int cur = 0:
      for (auto ch : s) {
        int &nxt = next[cur][int(ch - 'a')];
        if (!nxt) nxt = newnode();
         cnt[cur = nxt] += 1;
     }
   void build() {
     queue < int > q;
```

return

 $(0 \le i \text{ and } i < (int)z.size()) ? z[i] : 0; };$

```
q.push(0);
                                                                     auto add rep
    while (!q.empty()) {
                                                                           = [\&](bool left, int c, int l, int k1, int k2) {
      int cur = q.front();
                                                                             int L = \max(1, l - k2), R = \min(l - left, k1);
       q.pop();
       for (int i = 0; i < CNUM; ++i)
  if (next[cur][i])</pre>
                                                                        if (L > R) return;
                                                                        if (left)
           q.push(insertSAM(cur, i));
                                                                             rep[l].emplace\_back(sft + c - R, sft + c - L);
                                                                        else rep[l].emplace_back
                                                                             (sft + c - R - l + 1, sft + c - L - l + 1);
     vector < int > lc(tot);
    for (int i = 1; i < tot; ++i) ++lc[len[i]];
    partial_sum(ALL(lc), lc.begin());
                                                                     for (int cntr = 0; cntr < n; cntr++) {
                                                                       int 1, k1, k2;
    for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
                                                                        if (cntr < nu) {
                                                                          l = nu - cntr;
  void solve() {
                                                                          k1 = get_z(z1, nu - cntr);
    for (int i = tot - 2; i >= 0; --i)
                                                                          k2 = get_z(z2, nv + 1 + cntr);
       cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
                                                                        } else {
                                                                          l = cntr - nu + 1;
                                                                          k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
                                                                          k2 = get_z(z4, (cntr - nu) + 1);
     PalTree* [d7d2cf]
                                                                        if (k1 + k2 >= 1)
struct palindromic_tree {
                                                                          add_rep(cntr < nu, cntr, l, k1, k2);
  struct node {
    int next[26], fail, len;
                                                                  | \} // p \in [1, r] \Rightarrow s[p, p+i) = s[p+i, p+2i)
    int cnt, num; // cnt: appear times, num: number of
    // pal. suf. node(int l = 0) : fail(0), len(1), cnt(0), num(0) {
                                                                   6 Math
                                                                   6.1 ax+by=gcd(only exgcd *) [7b833d]
       for (int i = 0; i < 26; ++i) next[i] = 0;
    }
                                                                   pll exgcd(ll a, ll b)
  };
                                                                      if (b = 0) return pll(1, 0);
  vector<node> St;
                                                                      \hat{p} = a / b;
  vector<char> s;
                                                                     pll q = \operatorname{exgcd}(b, a \% b);
  int last , n;
                                                                     return pll(q.Y, q.X - q.Y * p);
  \begin{array}{ll} palindromic\_tree() : St(2), \ last(1), \ n(0) \ \{ \\ St[0].fail = 1, \ St[1].len = -1, \ s.pb(-1); \end{array}
                                                                   /* ax+by=res, let x be minimum non-negative
                                                                   g, p = gcd(a, b), exgcd(a, b) * res / g
                                                                   if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g) else: t = -(p.X / (b / g)) p += (b / g, -a / g) * t */
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
St.pb(0), St.pb(-1);
St[0].fail = 1, s.pb(-1);
                                                                   6.2 Floor and Ceil [692c04]
  inline int get_fail(int x) {
while (s[n - St[x].len - 1] != s[n])
                                                                   int floor (int a, int b)
                                                                   { return \hat{a} / \hat{b} - (a % \hat{b} && (a < 0) \hat{a} (b < 0)); }
      x = \hat{S}t[x].fail;
                                                                   int ceil(int a, int b)
                                                                   { return a / b + (a % b & (a < 0) \hat{} (b > 0)); }
    return x;
                                                                   6.3 Floor Enumeration [7cbcdf]
  inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
                                                                    / enumerating x = floor(n / i), [l, r]
                                                                   for (int l = 1, r; l \le n; l = r + 1) {
int x = n / l;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
       int now = SZ(St);
                                                                     r = n / x;
       St.pb(St[cur].len + 2);
      St[now].fail =
                                                                   6.4 Mod Min [9118e1]
         St[get_fail(St[cur].fail)].next[c];
                                                                   // \min\{k \mid 1 \le ((ak) \mod m) \le r\}, no solution -> -1 ll \max_{\min}(1l \ a, 1l \ m, 1l \ l, 1l \ r)  { if (a == 0) return l \ ? \ -1 : 0;
       St[cur].next[c] = now;
       St[now].num = St[St[now].fail].num + 1;
                                                                      if (11 \ k = (1 + a - 1) / a; \ k * a <= r)
    last = St[cur].next[c], ++St[last].cnt;
                                                                       return k;
  inline void count() { // counting cnt
                                                                      11 b = m / a, c = m \% a;
    auto i = St.rbegin();
                                                                     \quad \text{for } (; \ i \ != \ \operatorname{St.rend}(); \ +\!\!\!+\!\! i) \ \{
                                                                       return (1 + y * c + a - 1) / a + y * b;
      St[i->fail].cnt += i->cnt;
    }
                                                                   }
                                                                   6.5 Linear Mod Inverse [5a4cbf]
  inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
                                                                   inv[1] = 1;
                                                                   for (int i = 2; i
                                                                         \langle = N; ++i \rangle inv[i] = ((mod-mod/i)*inv[mod\%i])\%mod;
};
                                                                   6.6 Linear Filter Mu [ac2ac3]
5.10 Main Lorentz [615b8f]
                                                                   void getMu() {
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
                                                                     for (int i = 2; i \le n; ++i) {
                                                                       if (!flg[i]) p[++tot] = i, mu[i] = -1;
for (int j = 1; j <= tot && i * p[j] <= n; ++j) {
 flg[i * p[j]] = 1;
  if (i % p[i] = 0) {
  const int n = s.size();
  if (n = 1) return;
  const int nu = n / 2, nv = n - nu;
                                                                          if (i % p[j]
mu[i * p[j
  const string u = s.substr(0, nu), v = s.substr(nu),
                                                                                        = 0) \{
         ru(u.rbegin
                                                                                    p[j] = 0;
              (), u.rend()), rv(v.rbegin(), v.rend());
                                                                            break:
  main\_lorentz(u, sft), main\_lorentz(v, sft + nu)
  mu[i * p[j]] = -mu[i];
```

6.7 Gaussian integer gcd [763e59]

6.8 GaussElimination [6308be]

6.9 floor sum* [49de67]

6.10 Miller Rabin* [06308c]

6.11 Big number [6d475b]

```
template<typename T>
inline string to_string(const T& x) {
   stringstream ss;
   return ss<x,ss.str();
}
struct bigN:vector<ll>{
   const static int base=1000000000,width=log10(base);
   bool negative;
   bigN(const_iterator
        a,const_iterator b):vector<ll>(a,b) {}
```

```
bigN(string s){
  if(s.empty())return;
if(s[0]=='-')negative=1,s=s.substr(1);
  else negative=0;
  for (int i=int(s.size())-1; i>=0; i-=width) {
    11 t=0;
    for(int j=max(0,i-width+1);j \le i;++j)
       t=t*10+s[j]-'0';
    push_back(t);
  trim();
template<typename T>
  bigN(const T &x):bigN(to_string(x)){}
bigN(): negative(0){}
void trim(){
  while(size()&&!back())pop_back();
  if (empty()) negative=0;
void carry(int _base=base){
  for(size_t i=0;i<size();++i){
    if (at(i)>=0&&at(i)<_base)continue;
     if(i+1u=size())push\_back(0);
    int r=at(i)%_base;
    if(r<0)r=\_base;
    at(i+1)+=(at(i)-r)/\_base, at(i)=r;
int abscmp(const bigN &b)const{
  if(size()>b.size())return 1;
  if (size () < b. size () return -1;
  for (int i=int(size())-1;i>=0;--i){
     if(at(i)>b[i])return 1;
     if (at(i)<b[i]) return -1;
  return 0;
int cmp(const bigN &b)const{
  if (negative!=b.negative)return negative?-1:1;
  return negative?-abscmp(b):abscmp(b);
bool operator < (const bigN&b) const { return cmp(b) < 0;}
bool operator > (const bigN&b) const {return cmp(b) > 0;}
bool operator <= (const bigN&b) const {return cmp(b) <= 0;}
bool operator>=(const bigN\&b)const\{return cmp(b)>=0;\}
bool operator==(const bigN&b) const {return !cmp(b);}
bool operator!=(const bigN&b)const{return cmp(b)!=0;}
bigN abs()const{
  bigN res=*this;
  return res.negative=0, res;
bigN operator - () const {
  bigN res=*this;
  return res.negative=!negative, res.trim(), res;
bigN operator+(const bigN &b)const{
  if(negative)return -(-(*this)+(-b));
  if(b.negative)return *this-(-b);
  bigN res=*this;
  if (b. size()>size()) res. resize(b. size());
  for(size_t i=0; i < b. size(); ++i) res[i] +=b[i];
  return res.carry(),res.trim(),res;
bigN operator - (const bigN &b) const {
  if (negative) return -(-(*this)-(-b));
if (b. negative) return *this+(-b);
  if(abscmp(b)<0)return -(b-(*this));</pre>
  bigN res=*this;
  if (b. size()>size()) res. resize(b. size());
  for(size_t' i=0;i<b.size();++i)res[i]-=b[i];
  return res.carry(), res.trim(), res;
bigN operator*(const bigN &b)const{
  bigN res;
  {\tt res.negative} \small = {\tt negative!} \small = {\tt b.negative};
  res.resize(size()+b.size())
  for(size_t i=0; i < size(); ++i)
    for (size_t j=0; j< b. size();++j)
       if ((res[i+j]+=at(i)*b[j])>=base){
  res[i+j+l]+=res[i+j]/base;
         res\:[\:i+j]\% = base\:;
       \frac{1}{4}ak¥ carry · | ·
  return res.trim(),res;
bigN operator/(const bigN &b)const{
  int norm=base/(b.back()+1);
```

```
National Yang Ming Chiao Tung University FubukiMyWife
    bigN = abs()*norm;
    bigN y=b.abs()*norm;
    bigN q,r;
    q.resize(x.size());
    for (int i=int(x.size())-1;i>=0;--i){
      r=r*base+x[i];
      int s1=r.size()<=y.size()?0:r[y.size()];</pre>
      int s2=r.size()<y.size()?0:r[y.size()-1];
      int d=(ll(base)*s1+s2)/y.back();
      r=r-y*\dot{d};
      while (r.negative) r=r+y, --d;
      q[i]=d;
    q.negative=negative!=b.negative;
    return q.trim(),q;
  bigN operator%(const bigN &b)const{
    return *this-(*this/b)*b;
  friend istream& operator>>(istream &ss, bigN &b) {
    string s;
    return ss>>s, b=s, ss;
       ostream& operator << (ostream &ss, const bigN &b) {
    if (b.negative)ss<<'-';
ss<<(b.empty()?0:b.back());</pre>
    for (int i=int(b.size())-2;i>=0;--i)
      ss << setw (width) << set fill ('0') << b[i];
    return ss;
  template<typename T>
    operator T(){
      stringstream ss;
      ss<<*this;
      T res;
      return ss>>res,res;
};
6.12
      Fraction [4ab37a]
struct fraction {
  ll n, d;
  fraction
      (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
    ll\ t\ =\ \gcd(n\,,\ d)\,;
    n \neq t, d \neq t;
    if'(d < 0)' n = -n, d = -d;
```

```
struct fraction {
    ll n, d;
    fraction
        (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
        ll t = gcd(n, d);
        n /= t, d /= t;
        if (d < 0) n = -n, d = -d;
    }
    fraction operator-() const
    { return fraction(-n, d); }
    fraction operator+(const fraction &b) const
    { return fraction(n * b.d + b.n * d, d * b.d); }
    fraction operator-(const fraction &b) const
    { return fraction(n * b.d - b.n * d, d * b.d); }
    fraction operator*(const fraction &b) const
    { return fraction(n * b.n, d * b.d); }
    fraction operator/(const fraction &b) const
    { return fraction(n * b.d, d * b.n); }
    void print() {
        cout << n;
        if (d != 1) cout << "/" << d;
    }
};</pre>
```

6.13 Simultaneous Equations [a231be]

```
struct matrix { //m variables, n equations
  \label{eq:maxn} \textit{fraction} \ \ M[\textit{MAXN}] \left[ \textit{MAXN} + \ 1 \right], \ \ \textit{sol} \left[ \textit{MAXN} \right];
  int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;
       if (piv == m) continue;
            (int j = 0; j < n; +++j)  {
          if (i == j) continue;
          fraction\ tmp = -M[\,j\,]\,[\,piv\,]\ /\ M[\,i\,]\,[\,piv\,]\,;
          for (int k = 0; k < 0
                m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
       }
     int rank = 0;
     for (int i = 0; i < n; ++i) {
       int piv = 0;
       while (piv < m \& M[i][piv].n) ++piv;
```

6.14 Pollard Rho* [fdef9b]

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n = 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2
      = 0) return PollardRho(n / 2), ++cnt[2], void();
   11 x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
   while (true) {
    if (d != n & d != 1) {
       PollardRho(n / d);
       PollardRho(d);
       return;
    if (d = n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = \gcd(abs(x - y), n);
}
```

6.15 Simplex Algorithm [6b4566]

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
\begin{array}{l} \textbf{double} \ d \, [\text{MAXN}] \, [\text{MAXM}] \, , \ x \, [\text{MAXM}] \, ; \\ \textbf{int} \ ix \, [\text{MAXN} + \text{MAXM}] \, ; \ // \ !!! \ array \ all \ indexed \ from \ 0 \end{array}
 // \max\{cx\} \text{ subject to } \{Ax \le b, x > = 0\}
         n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
  // value = simplex(a, b, c, N, M);
double simplex(int n, int m){
         fill_n(d[n], m + 1, 0);
        fill_n(d[n+1], m+1, 0);
        iota(ix, ix + n + m, 0);
         int r = n, s = m - 1;
         for (int i = 0; i < n; ++i) {
               for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];
               d[i][m-1] = 1;

d[i][m] = b[i];
                if (d[r][m] > d[i][m]) r = i;
       copy_n(c, m - 1, d[n]);
       d[n + 1][m - 1] = -1;
       for (double dd;; ) {
                if (r < n) {
                       swap(ix[s], ix[r+m])
                       \begin{aligned} & \text{swap}(\text{IX}[s], \text{ IX}[\text{I} + \text{IM}]), \\ & \text{d}[\text{r}][s] = 1.0 / \text{d}[\text{r}][s]; \\ & \text{for (int } j = 0; \text{ } j <= \text{m}; ++j) \\ & \text{if } (j != s) \text{d}[\text{r}][j] \begin{subarray}{c} *= -\text{d}[\text{r}][s]; \\ *= -\text{d}[s][s]; \\ *= -\text{d}[s][s][s]; \\ *= -\text{d}[s][s][s]; \\ *= -\text{d}[s][s][s]; \\ *= -\text{d}[s][s][s][s]; \\ *= -\text{d
                       for (int i = 0; i \le n + 1; ++i) if (i != r) {
                             for (int j = 0; j <= m; ++j) if (j != s) d[i][j] += d[r][j] * d[i][s]; d[i][s] *= d[r][s];
                       }
                r = s = -1;
              for (int j = 0; j < m; ++j)

if (s < 0 \mid | ix[s] > ix[j]) {

if (d[n + 1][j] > eps \mid |

(d[n + 1][j] > -eps && d[n][j] > eps))
                if(s < 0) break;
               for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
                       if(r < 0 | |
                                       (dd = d[r][m]
                                                          / d[r][s] - d[i][m] / d[i][s]) < -eps | |
                                       (dd < eps \, \&\& \, ix \, [\, r \, + m] \, > \, ix \, [\, i \, + m] \, ) \, )
                               r = i;
                if (r < 0) return -1; // not bounded
        if (d[n + 1][m] < -eps) return -1; // not executable
       double ans = 0;
```

```
\begin{array}{ll} fill\_n\,(x\,,\,\,m,\,\,0)\,;\\ for\,\,(\,int\,\,\,i\,=\,m;\,\,i\,<\,\,\\ \end{array}
   n+m; \ ++i) \ \ \{ \ // \ \ the \ missing \ enumerated \ x[i]=0 if (ix[i] < m - 1){
       ans += d[i - m][m] * c[ix[i]];
       x[ix[i]] = d[i-m][m];
return ans;
```

6.15.1 Construction

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{T} y$ s.t. $A^{T} y \ge c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c$

 $\overline{\mathbf{x}}$ and $\overline{\mathbf{y}}$ are optimalified only if for all $i \in [1, n]$, either $\overline{x}_i = 0$ or $\sum_{j=1}^m A_{ji} \overline{y}_j = c_i$ holds and for all $i \in [1, m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij} \bar{x}_j = b_j$ holds.

```
1. In case of minimization, let c'_i = -c_i
```

- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- $3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j$
 - $\bar{\sum}_{1 \leq i \leq n}^{-} A_{ji} x_i \leq b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.16 chineseRemainder [a53b6d]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
  \begin{array}{l} ll \;\; g = \gcd{(m1,\;m2)}\,; \\ if \;\; ((x2 \; - \; x1) \;\% \; g) \;\; return \;\; \text{-1}; \; \text{// no sol} \end{array}
  m1 /= g; m2 /= g;
   pll p = exgcd(m1, m2);
  ll lcm = m1 * m2 * g;
ll res = p.first * (x^2 - x^1) * m1 + x1;
   // be careful with overflow
   return (res % lcm + lcm) % lcm;
```

6.17 Factorial without prime factor* [c324f3]

```
O(p^k + \log^2 n), pk = p^k
 11 prod [MAXP];
 ll fac_no_p(ll n, ll p, ll pk) {
   \operatorname{prod}[0] = 1;
   for (int i = 1; i <= pk; ++i)
if (i % p) prod[i] = prod[i - 1] * i % pk;
      else prod[i] = prod[i - 1];
   11 \text{ rt} = 1;
   for (; n; n /= p) {
     rt = rt * mpow(prod[pk], n / pk, pk) % pk;
rt = rt * prod[n % pk] % pk;
   return rt;
} // (n! without factor p) % p^k
```

6.18 PiCount* [cad6d4]

```
ll PrimeCount(ll n) { // n ~ 10^13 \Longrightarrow < 2s
  if (n \le 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector < int > smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i \le v; ++i) smalls [i] = (i + 1) / 2;
  for (int i = 0; i < s; ++i) { roughs [i] = 2 * i + 1;
    larges [i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p \le v; ++p) {
    if (smalls[p] > smalls[p - 1]) {
        int q = p * p;
       ++pc;
       if (1LL * q * q > n) break;
       skip[p] = 1;
       for (int i = q; i \le v; i += 2 * p) skip[i] = 1;
       int \dot{n}s = 0;
       for (int k = 0; k < s; ++k) {
         int i = roughs[k];
         if (skip[i]) continue;
11 d = 1LL * i * p;
         larges [ns] = larges [k] - (d \le v ? larges
               [smalls[d] - pc] : smalls[n / d]) + pc;
         roughs [ns++] = i;
       for (int j = v / p; j >= p; --j) {
         int c =
                smalls[j] - pc, e = min(j * p + p, v + 1);
```

```
16
           for (int i = j * p; i < e; ++i) smalls[i] -= c;
     }
   for (int k = 1; k < s; ++k) {
     const ll m = n / roughs [k];
ll t = larges [k] - (pc + k - for (int l = 1; l < k; ++l) {
        int p = roughs[1];
if (1LL * p * p > m) break;
t -= smalls[m / p] - (pc + 1 - 1);
     larges[0] = t;
   return larges[0];
}
6.19 Discrete Log* [da27bf]
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered\_map < int, int > p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {
     p[y] = i;

y = 1LL * y * x % m;

b = 1LL * b * x % m;
```

for (int i = 0; i < m + 10; i += kStep) { s = 1LL * s * b % m; if (p.find(s) != p.end()) return i + kStep - p[s]; return -1; int DiscreteLog(int x, int y, int m) { if (m == 1) return 0; int s = 1;for (int i = 0; i < 100; ++i) { if (s == y) return i; s = 1LL * s * x % m;

int p = 100 + DiscreteLog(s, x, y, m);

6.20 Berlekamp Massey [3eb6fa]

if (fpow(x, p, m) != y) return -1;

if (s = y) return 100;

return p;

```
template <typename T>
  vector<T> BerlekampMassey(const vector<T> &output) {
                    vector < T > d(SZ(output) + 1), me, he;
                   for (int f = 0, i = 1; i <= SZ(output); ++i) {
  for (int j = 0; j < SZ(me); ++j)
   d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
                                       if (me.empty()) {
                                                    me. resize (f = i);
                                                       continue;
                                       vector < T > o(i - f - 1);
                                   T k = -d[i] / d[f]; o.pb(-k);
for (T x : he) o.pb(x * k);
                                    for (i \times i \times i) o. (i \times i)
                                   me = o;
                   return me;
```

6.21 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679\ 999983\ 1097774749\ 1076767633\ 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 \ 98789101 \ 987777733 \ 999991921 \ 1010101333
     1010102101 \ 1000000000039 \ 100000000000037
     2305843009213693951 \ \ 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

```
Estimation
 n | 2345 6 7 8 9 20 30 40 50 100
p(n) 23571115223062756044e42e52e8
 n |1001e31e6 1e9 1e12 1e15 1e18
d(i) 12 32 2401344672026880103680
 n \mid 12345678
                               9
                                       10
                                           11 12 13 14 15
\binom{2n}{n} \, 2\, 6\, 20\, 70\, 252\, 924\, 3432\, 12870\, 48620\, 184756\, 7e5\, 2e6\, 1e7\, 4e7\, 1.5e8
n 2 3 4 5 6 7 8 9
                                10 11 12 13
B_n 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7
```

6.23 General Purpose Numbers

Bernoullinumbers $B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$ $\sum_{j=0}^{m} \binom{m+1}{j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into $\operatorname{exactly} k \operatorname{groups}$.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers 1

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. kj:ss.t. $\pi(j) > \pi(j+1), k+1j$:ss.t. $\pi(j) \ge j$, kj:ss.t. $\pi(j) > j$.

E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.24 Tips for Generating Functions

Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$

```
-A(rx) \Rightarrow r^n a_n
-A(x)+B(x) \Rightarrow a_n+b_n
-A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}
```

 $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$

 $-xA(x)' \Rightarrow na_n$ $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$

• Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$

 $-A(x)+B(x) \Rightarrow a_n+b_n$

 $-A^{(k)}(x) \Rightarrow a_{n+k}$ $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$

 $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$

 $-xA(x) \Rightarrow na_n$

 ${\bf Special \, Generating \, Function}$

 $- (1+x)^n = \sum_{i \ge 0} \binom{n}{i} x^i$

$rac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{i}{n-1} x^i$ $extbf{Polynomial}$

Fast Fourier Transform [5e2ea2]

```
const int maxn = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
   for (int i = 0; i <= maxn</pre>
         ; ++i) \text{ omega[i]} = \exp(i * 2 * pi / maxn * I);
void bin(vector<cplx> &a, int n) {
    int lg;
    for (lg = 0; (1 << lg) < n; ++lg); --lg;
    vector < cplx > tmp(n);
    for (int i = 0; i < n; ++i) {
         int to = 0;
         for (int j = 0; (1 << j) <
            n; ++j) to |= (((i >> j) \& 1) << (lg - j));
        tmp[to] = a[i];
    for (int i = 0; i < n; ++i) a[i] = tmp[i];
void fft(vector<cplx> &a, int n) {
    bin(a, n);
    for (int step = 2; step \ll n; step \ll 1) {
        int to = step >> 1;
for (int i = 0; i < n; i += step) {
             for (int k = 0; k < to; ++k) {
                 cplx x = a[i]
                     + to + k] * omega[maxn / step * k];
                 a[i + to + k] = a[i + k] - x;
```

```
a[i + k] += x;
             }
        }
    }
}
void ifft(vector<cplx> &a, int n) {
    fft(a, n);
    reverse(a.begin() + 1, a.end());
    for (int i = 0; i < n; ++i) a[i] /= n;
vector<int> multiply(const vector<
    int>&a, const vector<int>&b, bool trim = false) {
    int d = 1;
    while
        (d < \max(a.size(), b.size())) \ d <\!\!<= 1; \ d <\!\!<= 1;
    vector < cplx > pa(d), pb(d);
    for (int i
         = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
    for (int i
         = 0; i < b.size(); ++i) pb[i] = cplx(b[i], 0);
    fft (pa, d); fft (pb, d);
    for (int i = 0; i < d; ++i) pa[i] *= pb[i];
    ifft (pa, d);
    vector<int> r(d);
    for (int
         i = 0; i < d; ++i) r[i] = round(pa[i].real());
    if (trim)
        while (r.size() \& r.back() = 0) r.pop_back();
    return r;
                               Root
 Prime
            Root
                   Prime
                   167772161
 7681
            17
                               3
 12289
                   104857601
            11
                               3
 40961
            3
                   985661441
                               3
 65537
            3
                   998244353
                               3
 786433
            10
                   1107296257
                               10
 5767169
                   2013265921
            3
                               31
                   2810183681
 7340033
            3
                               11
 23068673
            3
                   2885681153
                               3
            3
                   605028353
 469762049
7.2 Number Theory Transform* [7d51db]
```

```
vector<int> omega;
 void Init() {
   omega. resize(kN + 1);
   long long x = \text{fpow}(kRoot, (Mod - 1) / kN);
   omega[0] = 1;
   for (int i = 1; i <= kN; ++i) {
  omega[i] = 1LL * omega[i - 1] * x % kMod;
void Transform (vector < int > &v, int n) {
   BitReverse(v, n);
for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
     for (int i = 0; i < n; i += s) {
        for (int k = 0; k < z; ++k) {
          v[i + k + z] = (v[i + k] + kMod' - x) \% kMod;
(v[i + k] += x) \% = kMod;
     }
 void InverseTransform(vector<int> &v, int n) {
   Transform \left( v \, , \ n \right);
   for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
   const int kInv = fpow(n, kMod - 2);
   for (int i
         = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;
}
```

7.3 Fast Walsh Transform* [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
   for (int L = 2; L \le n; L \le 1)
      for (int i = 0; i < n; i += L)
for (int j = i; j < i + (L >> 1); ++j)
a[j + (L >> 1)] += a[j] * op;
```

```
National Yang Ming Chiao Tung University FubukiMyWife
                                                                                          \begin{array}{l} X = rhs.Mul(Q)\,,\; Y = *this\,;\\ fi\,(0\,,\; n())\;\; if\;\; ((Y[\,i\,]\; -=\; X[\,i\,])\,<\,0)\;\; Y[\,i\,]\; +=\, P;\\ return\;\; \{Q,\; Y.\, is\, z\, (max(\,1\,,\; rhs.\, n()\;\; -\; 1))\,\}; \end{array}
const int N = 21;
int f
     N = [1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
                                                                                       Poly Dx() const {
    subset\_convolution(int *a, int *b, int *c, int L) \ \{ \ / \ c_k = \sum_{i=1}^{n} \{i \ | \ j = k, \ i \ \& \ j = 0 \} \ a_i * b_j \ \} \ \}
                                                                                          Poly ret(n() - 1);
                                                                                          fi(0,
  int n = 1 \ll L;
  for (int i = 1; i < n; +++i)
     ct[i] = ct[i \& (i - 1)] + 1;
  for (int i = 0; i < n; ++i)
                                                                                       Poly Sx() const {
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
                                                                                          Poly ret(n() + 1);
  for (int i = 0; i \le L; ++i)
                                                                                          fi(0, n())
     fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
                                                                                          return ret;
     for (int j = 0; j <= i; ++j)
for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];
                                                                                       Polv
  for (int i = 0; i \le L; ++i)
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
     c[i] = h[ct[i]][i];
                                                                                          if (!m) return { };
7.4 Polynomial Operation [869cb1]
                                                                                          // fi(2, m *
\begin{array}{ll} fi(s,\ n) \ for \ (int\ i=(int)(s); \ i<(int)(n); \ +\!\!+\!\!i) \\ template<int\ MAXN, \ ll\ P, \ ll\ RT>//\ MAXN=2^k \end{array}
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<ll>>::vector;
  static NTTMAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
                                                                                          vector<ll> y(m);
  Poly(const Poly &p, int m) : vector<ll>(m) {
     copy_n(p.data(), min(p.n(), m), data());
                                                                                          return y;
  Poly& irev()
                                                                                        static vector < Poly>
  { return reverse(data(), data() + n()), *this; } Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) \{ // n() = rhs.n() \}
     fi(0, n()) if
  (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
  Poly& imul(ll k) {
     fi(0, n()) (*this)[i] = (*this)[i] * k % P;
return *this;
                                                                                       vector
  Poly Mul(const Poly &rhs) const {
     int m = 1;
     while (m < n() + rhs.n() - 1) m <<= 1;
     Poly X(*this, m), Y(rhs, m);
     \begin{array}{l} {\rm ntt}\,(X.\,{\rm data}\,()\,,\,\,m)\,,\,\,\,{\rm ntt}\,(Y.\,{\rm data}\,()\,,\,\,m)\,;\\ {\rm fi}\,(\,0\,,\,m)\,\,X[\,{\rm i}\,]\,=\,X[\,{\rm i}\,]\,\,\,{}^*\,\,Y[\,{\rm i}\,]\,\,\%\,\,P; \end{array}
     ntt(X.data(), m, true);
     return X.isz(n() + rhs.n() - 1);
                                                                                          for (int i = m -
  Poly Inv() const \{ // (*this)[0] != 0, 1e5/95ms \}
     if (n() = 1) return \{ntt.minv((*this)[0])\};
                                                                                          return down[1];
     while (m < n() * 2) m <<= 1;
Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
     Poly Y(*this, m);
     ntt(Xi.data(), m), ntt(Y.data(), m);
     fi (0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;
```

ntt(Xi.data(), m, true);

const $\{ // \text{ Jacobi}((*this)[0], P) = 1, 1e5/235ms \}$

= 1) return {QuadraticResidue((*this)[0], P)};

 $X = Poly(*this \,, \ (n() \,+\, 1) \,\,/\,\, 2).\,Sqrt()\,.\,isz(n())\,;$

X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);

(const Poly &rhs) const { // (rhs.)back() != 0

if (n() < rhs.n()) return $\{\{0\}, *this\};$

Poly Q = Y.Mul(X.Inv()).isz(m).irev();

const int m = n() - rhs.n() + 1;

Poly X(rhs); X. irev(). isz(m);

Poly Y(*this); Y. irev(). isz(m);

return Xi.isz(n());

pair<Poly, Poly> DivMod

Poly Sqrt()

if (n()

Poly

```
ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
      return ret.isz(max(1, ret.n()));
                   ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
               _tmul(int nn, const Poly &rhs) const {
     Poly Y = Mul(rhs).isz(n() + nn -
     return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const
             vector<ll> &x, const vector<Poly> &up) const {
      const int m = (int)x.size();
      vector < Poly > down (m * 2);
      // \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};
                   2) \ \operatorname{down}[\ i\ ] \ = \ \operatorname{down}[\ i\ /\ 2\ ] \ . \ \operatorname{DivMod}(\operatorname{up}[\ i\ ]) \ . \ \operatorname{second};
     down[1] = Poly(up[1])
     \begin{array}{ll} & ... \\ ... \\ ... \\ irev(). \\ ... \\ isz(n()). \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... \\ ... 
                                             1]._tmul(up[i].n() - 1, down[i / 2]);
      fi(0, m) y[i] = down[m + i][0];
                                                         _{\text{tree1}(\text{const} \text{ vector} < ll > \&x)}  {
     const int m = (int)x.size();
vector<Poly> up(m * 2);
     fi(0, m) up[m + i] = {(x[i] ? P - x[i] : 0), 1};

for (int i = m - 1; i

> 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
           <ll> Eval(const vector <ll> &x) const { // 1e5, 1s}
     auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate (const vector
          <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
     const int m = (int)x.size();
     vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
fi(0, m) down[m + i] = {z[i]};
                 `1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
     return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] == 0, 1e5/360ms
     if (n() = 1) return \{1\};

Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());

Poly Y = X.Ln(); Y[0] = P - 1;
      fi(0, n())
                   if'((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;
     return X.Mul(Y).isz(n());
  V/M := P(P - 1). If k >= M, k := k \% M + M.
Poly Pow(ll k) const {
     int nz = 0;
     while (nz < n() \&\& !(*this)[nz]) ++nz;
if (nz * min(k, (ll)n()) >= n()) return Poly(n());
if (!k) return Poly(Poly \{1\}, n());
      Poly X(data() + nz, data() + nz + n() - nz * k);
     const ll c = ntt.mpow(X[0], k % (P - 1));
return X.Ln().imul
                 (k % P).Exp().imul(c).irev().isz(n()).irev();
static 11
           LinearRecursion(const vector<ll> &a, const vector
           <ll> &coef , ll n) { // a_n = \sum c_j a_(n-j)
      const int k = (int)a.size();
```

```
assert((int)coef.size() == k + 1);
Poly C(k + 1), W(Poly {1}, k), M = {0, 1};
fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
C[k] = 1;
while (n) {
    if (n % 2) W = W.Mul(M).DivMod(C).second;
    n /= 2, M = M.Mul(M).DivMod(C).second;
}
ll ret = 0;
fi(0, k) ret = (ret + W[i] * a[i]) % P;
return ret;
}
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template ⇒ decltype(Poly_t::ntt) Poly_t::ntt = {};
```

7.5 Value Polynomial [96cde9]

```
{\color{red} \textbf{struct}} \hspace{0.1cm} \textbf{Poly} \hspace{0.1cm} \{
  mint base; // f(x) = poly[x - base]
  vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) \{\}
  mint get_val(const mint &x) {
    if (x >= base & x < base + SZ(poly))
      return poly[x - base];
    mint rt = 0;
    for (int i = 0; i < SZ(poly); ++i)
  rt += poly[i] * ifac[i] * inegfac
      [SZ(poly) - 1 - i] * lmul[i] * rmul[i];</pre>
    return rt;
  return;
    mint nw = get_val(base + SZ(poly));
    poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)
       poly[i] += poly[i - 1];
};
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P) = 0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

8 Geometry

8.1 Basic [8c3633]

```
double a, b, c, o;
   P pa, pb;
   \begin{array}{l} P \; pa, \; pb, \\ L() \; : \; a(0), \; b(0), \; c(0), \; o(0), \; pa(), \; pb() \; \{\} \\ L(P \; pa, \; P \; pb) \; : \; a(pa.y \; - \; pb.y), \; b(pb.x \; - \; pa.x \\ \quad \  \  \, ), \; c(pa \; \hat{} \; pb), \; o(atan2(-a, \; b)), \; pa(pa), \; pb(pb) \; \{\} \end{array}
   P project(P p) { return pa + (pb - pa).unit
    () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs()) * (pb - pa).abs()); }
};
bool SegmentIntersect (P p1, P p2, P p3, P p4) {
   if (\max(p1.x, p2.x) < \min(p3.x, p4.x)
            \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
    if (\max(p1.y, p2.y) < \min(p3.y, p4.y) | |
            \max(\texttt{p3.y}, \texttt{p4.y}) < \min(\texttt{p1.y}, \texttt{p2.y})) \text{ return false};
   return sign((p3 - p1) ^{\circ} (p4 - p1)) * sign((p3 - p2) ^{\circ} (p4 - p2)) <= 0 && sign((p1 - p3) ^{\circ}
                (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
       (L x, L y) { return same(x.a * y.b, x.b * y.a); }
       8.2 KD Tree [375ca2]
namespace kdt {
int root, lc [maxn],
         rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
    if (l = r) return -1;
   function < bool (const point &, const point
            &>> f = [dep](const point &a, const point &b) {
       if (dep \& 1) return a.x < b.x;
       else return a.y < b.y;
   int m = (1 + r) >> 1;
   nth\_element(p + 1, p + m, p + r, f);
   \begin{array}{l} \operatorname{xl}\left[m\right] = \operatorname{xr}\left[m\right] = \operatorname{p}\left[m\right].\,\mathrm{x}\,;\\ \operatorname{yl}\left[m\right] = \operatorname{yr}\left[m\right] = \operatorname{p}\left[m\right].\,\mathrm{y}\,; \end{array}
   lc[m] = build(l, m, dep + 1);
    if (~lc[m]) {
       xl[m] = min(xl[m], xl[lc[m]]);
        \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{lc}\left[m\right]\right]\right); \end{array} 
       yr[m] = max(yr[m], yr[lc[m]]);
   rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
       xl[m] = min(xl[m], xl[rc[m]]);
       \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{rc}[m]]);
       yl[m] = min(yl[m], yl[rc[m]]);

yr[m] = max(yr[m], yr[rc[m]]);
   return m:
bool bound(const point &q, int o, long long d) {
   double ds = sqrt(d + 1.0);
     if \ (q.x < xl [o] - ds \ | \ | \ q.x > xr [o] + ds \ | \ | \\
                  yl[o] - ds \mid | q.y > yr[o] + ds) return false;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
       const point &q, long long &d, int o, int dep = 0) {
    if (!bound(q, o, d)) return;
   long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1)
       && q.x < p[o].x || !(dep & 1) && q.y < p[o].y) { if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
       if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
   } else {
       if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
       if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
```

```
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
}
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
}}

8.3 Sector Area [c41fb7]

// calc area of sector which include a b
</pre>
```

8.4 Half Plane Intersection [f7274e]

```
bool jizz(L l1,L l2,L l3){
  P p=Intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;
bool cmp(const L &a, const L &b){
  return same(
       a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
vector<L> pls(1,ls[0]);
  for (int i=0; i<(int) ls. size ();++i) if (!
       same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
  deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c
    ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
  for (int i=2; i < (int) pls. size(); ++i)
    \begin{array}{l} meow(i,dq.back(),dq[dq.size()-2])dq.pop\_back();\\ meow(i,dq[0],dq[1])dq.pop\_front(); \end{array}
    dq.push_back(i);
  }
  meow(dq
        .front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop\_front();\\
  if (dq. size ()<3u) return vector
       <P>(); // no solution or solution is not a convex
  vector<P> rt;
  for (int i=0; i<(int)dq. size();++i)rt.push_back
       (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
  return rt:
}
```

8.5 Rotating Sweep Line [0411f0]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
  int n=int(ps.size());
  vector < int > id(n), pos(n);
  vector<pair<int, int>>> line(n*(n-1)/2);
  int m=-1:
  for(int i=0;i< n;++i)for
       (int j=i+1; j< n; ++j) line[++m] = make_pair(i,j); ++m;
  sort(line.begin(),line.end(),[&](const
       pair<int, int> &a, const pair<int, int> &b)->bool{
      if (ps
          [a.first].first = ps[a.second].first)return 0;
      if (ps
          [b.first].first=ps[b.second].first)return 1;
      return (double
           )(ps[a.first].second-ps[a.second].second)/(ps
           [a. first]. first - ps[a. second]. first) < (double)
           ) (ps[b.first].second-ps[b.second].second
           )/(ps[b.first].first-ps[b.second].first);
      });
  for (int i=0; i< n; ++i) id [i]=i;
  sort(id.begin(),id.end(),[&](const
       int &a, const int &b) { return ps[a] < ps[b]; });
  for (int i=0; i< n; ++i) pos[id[i]]=i;
  for (int i=0; i \leqslant m++i){
    auto l=line[i];
```

```
// meow
    tie (pos[1.first],pos[1.second],
        id [pos[1.first]],id [pos[1.second]])=make_tuple
        (pos[1.second],pos[1.first],1.second,1.first);
}
```

8.6 Triangle Center [4e8ee9]

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y)
  double bx = (c.x + b.x) /
  double by = (c.y + b.y) / 2;
double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by
- ay)) / (sin(a1) * cos(a2) - sin(a2) * cos(a1));
return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point\ TriangleOrthoCenter(Point\ a,\ Point\ b,\ Point\ c)\ \{
  return TriangleMassCenter(a, b
       , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);
  res.x = (
      la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (
    la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

8.7 Polygon Center [ee6ff0]

```
| Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

8.8 Maximum Triangle [3a6d38]

```
{\color{red} \textbf{double}} \ \ \textbf{ConvexHullMaxTriangleArea}
                                              (Point p[], int res[], int chnum) {
                         double area = 0, tmp;
                        res[chnum] = res[0];
                        for (int i = 0, j = 1, k = 2; i < chnum; i++) {
                                              while (fabs(Cross(p[
                                                                                      \begin{array}{lll} res[j]] & -p[res[i]] \; , \; p[res[(k+1) \; \% \; chnum]] \; -p[res[i]])) > fabs(Cross(p[res[j]] \; -p[res[i]] \; , \\ p[res[k]] \; -p[res[i]]))) \; k = (k+1) \; \% \; chnum; \\ \end{array} 
                                          tmp = fabs (Cross (
                                                                                         p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                              if (tmp > area) area = tmp;
                                              while (fabs(Cross(p[
                                                                                          \begin{array}{l} \operatorname{res}\left[\left(j+1\right)\%\operatorname{chnum}\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \ \operatorname{p}\left[\operatorname{res}\left[k\right]\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \\ \operatorname{p}\left[\operatorname{res}\left[i\right]\right]\right) > \operatorname{fabs}\left(\operatorname{Cross}\left(\operatorname{p}\left[\operatorname{res}\left[j\right]\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \\ \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right] - \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right] - \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right], \\ \operatorname{chnum}\left[\operatorname
                                                                                                p[res[k]] - p[res[i]]))) j = (j + 1) % chnum;
                                            tmp = fabs (Cross (
                                                                                         p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                              if (tmp > area) area = tmp;
                      return area / 2;
```

8.9 Point in Polygon [0a9a66]

```
int pip(vector<P> ps, P p) {
  int c = 0;
```

```
for (int i = 0; i < ps. size(); ++i)
                                                                         int j = (i + 1) \% 3;
    int a = i, b = (i + 1) \% ps.size();
                                                                         double o = atan2
                                                                         \begin{array}{lll} (ps\,[\,i\,].\,y,\ ps\,[\,i\,].\,x) & -\ atan2\,(ps\,[\,j\,].\,y,\ ps\,[\,j\,].\,x)\,;\\ if\ (o>=\ pi\,)\ o=o\ -\ 2\ *\ pi\,; \end{array}
    L l(ps[a], ps[b]);
                                                                         if (o >= pi) o = o - 2 * pi;
if (o <= -pi) o = o + 2 * pi;
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
                                                                         ans += AreaOfCircleTriangle
    if (same(ps[
                                                                              (ps[i]\,,\ ps[j]\,,\ r)\ *\ (o>=0\ ?\ 1\ :\ -1);
         a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
     if (ps[a].y > ps[b].y) swap(a, b);
     if (ps[a].y \le p.y \& p.y \le
                                                                       return abs(ans);
         \begin{array}{l} ps[b].y & & ps[a].x + (ps[b].x - ps[a].x \\ ) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) + c; \end{array}
                                                                    8.11
                                                                            Tangent of Circles and Points to Circle
  return (c & 1) * 2;
                                                                    vector<L> tangent(C a, C b) {
                                                                    #define Pij \
8.10 Circle [466c44]
                                                                      P \ i \ = \ (b.c \ - \ a.c) \ . \ unit \ () \ * \ a.r \ , \ j \ = \ P(\,i \ .y \ , \ -i \ .x) \ ; \ \backslash
struct C {
                                                                       z.emplace\_back(a.c + i, a.c + i + j);
  Р с;
                                                                    #define deo(I,J) \
  double r;
                                                                       double d = (a)
  C(P \ c = P(0, 0), \ double \ r = 0) : c(c), \ r(r) \ \{\}
                                                                           .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
                                                                      vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
                                                                       if (a.r < b.r) swap(a, b);
  vector<P> p;
                                                                       vector<L> z;
  if (same(a.r + b.r,
                                                                       if ((a.c - b.c).abs() + b.r < a.r) return z
        d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
                                                                       else if (same((a.c - b.c).abs() + b.r, a.r)) \{ Pij; \}
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
                                                                       else {
    double o = acos
                                                                         \operatorname{deo}(-,+);
         \left(\left(sq(a.r) \,+\, sq(d) \,-\, sq(b.r)\right) \,\,/\,\, \left(2 \,\,{}^*\,\, a.r \,\,{}^*\,\, d\right)\right);
                                                                         if (same(d, a.r + b.r)) \{ Pij; \}
else if (d > a.r + b.r) \{ deo(+,-); \}
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.rot(o) * a.r);
    p.push_back(a.c + i.rot(-o) * a.r);
                                                                       return z;
                                                                    }
  return p;
                                                                    vector<L> tangent(C c, P p) {
double IntersectArea(C a, C b) {
                                                                       vector < L > z;
  if (a.r > b.r) swap(a, b);
                                                                       \frac{double \ d = (p - c.c).abs();}{}
  double d = (a.c - b.c).abs();
                                                                       if(same(d, c.r)) {
  if (d \ge a.r + b.r - eps) return 0;
                                                                         P i = (p - c.c).rot(pi / 2);
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
                                                                         z.emplace\_back(p, p + i);
  double p = a\cos
                                                                       else if (d > c.r) 
       ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
                                                                         double o = acos(c.r / d);
  double q = acos
                                                                         P i = (p - c.c) \cdot unit'
  (), j = i.rot(o) * c.r, k = i.rot(-o) * c.r;
                                                                         z.emplace\_back(c.c + j, p);
                                                                         z.emplace\_back(c.c + k, p);
// remove second
level if to get points for line (defalut: segment) vector<P> CircleCrossLine(P a, P b, P o, double r) {
                                                                       return z;
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
                                                                    8.12 Area of Union of Circles [490636]
  double C = sq(a.x - o.x)
                                                                    vector<pair<double, double >>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
       ) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
                                                                       vector<pair<double, double>>> res;
  if (d \ge -eps)
    d = \max(0., d);
                                                                       if (same(a.r + b.i, u_{jj}), else if <math>(d \le abs(a.r - b.r) + eps) \{
                                                                       if (same(a.r + b.r, d))
    double i = (-B - \operatorname{sqrt}(d)) / (2 * A);
    double j = (-B + sqrt(d)) / (2 * A);
if (i - 1.0 \le eps \&\& i >=
                                                                       if (a.r < b.r) res.emplace_back(0, 2) else if (d < abs(a.r + b.r) - eps)
                                                                         -eps) t.emplace\_back(a.x + i * x, a.y + i * y);
     if (j - 1.0 <= eps && j >=
                                                                         if (z < 0) z += 2 * pi;
          -eps) t.emplace_back(a.x + j * x, a.y + j * y);
                                                                         double 1 = z - o, r = z + o;

if (1 < 0) 1 += 2 * pi;

if (r > 2 * pi) r -= 2 * pi;
  return t:
                                                                         if (l > r) res.emplace_back (l, 2 * pi), res.emplace_back(0, r);
// calc area
     intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
                                                                         else res.emplace_back(l, r);
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
                                                                       return res;
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
                                                                    double CircleUnionArea
    return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
                                                                         (\text{vector} < C > c) \{ // \text{ circle should be identical } 
                                                                       int n = c.size();
  if (inb) return
                                                                       double a = 0, w;
        SectorArea(p[0], a, r) + abs(p[0] ^ b)
                                                                       for (int i = 0; w = 0, i < n; ++i) {
     (p.size() = 2u) return SectorArea (a, p[0], r)
                                                                         vector<pair<double, double>>> s = {{2 * pi, 9}}, z; for (int j = 0; j < n; +j) if (i != j) {
        + SectorArea(p[1], b, r) + abs(p[0]
                                                     p[1]) / 2;
                                                                           z = CoverSegment(c[i], c[j]);
  else return SectorArea(a, b, r);
                                                                           for (auto &e : z) s.push_back(e);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
                                                                         sort(s.begin(), s.end());
                                                                         auto F = [\&] (double t) { return c[i].r * (c[i].r *
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
                                                                               t + c[i].c.x * sin(t) - c[i].c.y * cos(t)); };
```

```
for (auto &e : s) {
    if (e.first > w) a += F(e.first) - F(w);
    w = max(w, e.second);
    }
} return a * 0.5;

8.13 Minimun Distance of 2 Polygons [7eb8bb]
```

```
// p, q is convex double TwoConvexHullMinDist
      (Point P[], Point Q[], int n, int m) {
   int YMinP = 0, YMaxQ = 0;
   for (i =
          0; i < n; ++i) if(P[i].y < P[YMinP].y) YMinP = i;
   for (i =
          0; i < m; ++i) if(Q[i].y > Q[YMaxQ].y) YMaxQ = i;
   P[\, n\,] \; = P[\, 0\,] \; , \; \, Q[m] \; = Q[\, 0\,] \, ;
   for (int i = 0; i < n; ++i) {
      while (tmp = Cross(
             \begin{array}{l} \text{Q[YMaxQ + 1] - P[YMinP + 1], P[YMinP] - P[YMinP + 1])} > \text{Cross}\left(\text{Q[YMaxQ] - P[YMinP + 1], P[YMinP} \right) \end{array} 
      if (tmp < 0) ans = min(ans, PointToSegDist (P[YMinP], P[YMinP + 1]), Q[YMaxQ]));
       \begin{array}{l} \textbf{else} \ \ ans = \min(ans \,, \ TwoSegMinDist(P[ \\ YMinP] \,, \ P[YMinP \,+ \,\, 1] \,, \ Q[YMaxQ] \,, \ Q[YMaxQ \,+ \,\, 1])) \,; \end{array} 
      YMinP = (YMinP + 1) \% n;
   }
   return ans;
```

8.14 2D Convex Hull [65eaab]

```
bool operator < (const P &a, const P &b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator > (const P &a, const P &b) {
  #define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  \mathtt{sort}\,(\,\mathtt{ps.begin}\,(\,)\,\,,\,\,\,\mathtt{ps.end}\,(\,)\,\,,\,\,\,[\,\&\,]\,\,\,(\,\mathtt{P}\,\,\mathtt{a}\,,\,\,\,\mathtt{P}\,\,\mathtt{b}\,)\,\,\,\{\,\,\,\mathtt{return}\,\,
        same(a.x, b.x) ? a.y < b.y : a.x < b.x; });
    or (int i = 0; i < ps.size(); ++i) {
while (p.size() >= 2 \&\& crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t & crx(p[p.size() -
          [2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
 p.\,pop\_back\,(\,)\;;
  return p;
int sgn(double
      x) \{ return same(x, 0) ? 0 : x > 0 ? 1 : -1; \}
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector < P > p, u, d;
  CH() {}
 CH(vector < P > ps) : p(ps) {
    n = ps.size();
    rotate (p. begin
         ()\;,\;\min\_{\rm element}(p.\,begin\,()\;,\;p.\,end\,()\,)\;,\;p.\,end\,()\,)\;;
    auto t = max_element(p.begin(), p.end());
d = vector<P>(p.begin(), next(t));
    u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find (vector <P> &v, P d) {
    int l = 0, r = v.size();
     while (1 + 5 < r) {
```

```
else l = L;
     int x = 1;
     for (int i = l +
            1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
     return x;
   int findFarest(P v) {
     if (v.y > 0 | v.y = 0 \& v.x > 0) return
            ((int)d.size() - 1 + find(u, v)) % p.size();
     return find(d, v);
     get(int 1, int r, Pa, Pb) {
     int s = sgn(crx(a, b, p[1 \% n]));
     while (l + 1 < r) {
        int m = (l + r) >> 1;
        if (\operatorname{sgn}(\operatorname{crx}(a, b, p[m \% n])) == s) l = m;
        else r = m;
     return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
   vector <P> getLineIntersect(P a, P b) {
     int X = \text{findFarest}((b - a).\text{rot}(pi / 2));

int Y = \text{findFarest}((a - b).\text{rot}(pi / 2));

if (X > Y).\text{swap}(X, Y);
     if (sgn
            \begin{array}{l} (\operatorname{crx}(a,\ b,\ p[X])) \ * \ \operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[Y])) < 0) \\ \operatorname{return} \ \{ \operatorname{get}(X,\ Y,\ a,\ b) \,,\ \operatorname{get}(Y,\ X+n,\ a,\ b) \}; \end{array} 
     return {}; // tangent case falls here
   void update_tangent(P q, int i, int &a, int &b) {
     \begin{array}{ll} if & (sgn(crx(q,\ p[a]\,,\ p[i])) > 0) \ a = i\,; \\ if & (sgn(crx(q,\ p[b],\ p[i])) < 0) \ b = i\,; \end{array}
   void bs(int 1, int r, Pq, int &a, int &b) {
     if (l = r) return;
     update_tangent(q, 1 % n, a, b);
     int s = sgn(crx(q, p[1 \% n], p[(1 + 1) \% n]));

while (1 + 1 < r) {
   int m = (1 + r) >> 1;
        if (sgn(crx
              (q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
        else r = m;
     update_tangent(q, r % n, a, b);
   int x = 1;
  for (int i = l)
         + 1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
  return x;
int findFarest(P v) {
   if (v.y > 0 \mid v.y = 0 \& v.x > 0) return
         ((int)d.size() - 1 + find(u, v)) \% p.size();
   return find(d, v);
  get(int 1, int r, Pa, Pb) {
  int s = sgn(crx(a, b, p[l \% n]));
   while (l + 1 < r) {
     int m = (l + r) >> 1;
      \mbox{if } (sgn(crx(a,\ b,\ p[m\ \%\ n])) == s) \ l = m; \\
     else r = m;
  return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
vector <P> getIS (Pa, Pb) {
   \begin{array}{ll} \text{int } X = findFarest((b-a).spin(pi / 2));\\ \text{int } Y = findFarest((a-b).spin(pi / 2)); \end{array} 
   if (X > Y) swap(X, Y)
   if (\operatorname{sgn}(\operatorname{crx}(a, b, p[X])) * \operatorname{sgn}(\operatorname{crx}(a, b, p[Y])) <
        0) return \{ get(X, Y, a, b), get(Y, X + n, a, b) \};
  return { };
void update_tangent(P q, int i, int &a, int &b) { if (sgn(crx(q, p[a], p[i])) > 0) a = i;
  if (sgn(crx(q, p[b], p[i])) < 0) b = i;
void bs(int l, int r, Pq, int &a, int &b) {
  if (l == r) return;
  update_tangent(q, 1 % n, a, b);
    int \ s = sgn(crx(q, \ p[l \ \% \ n], \ p[(l + 1) \ \% \ n])); 
   while (1 + 1 < r)
     int m = (1 + r) >> 1;
```

```
if (sgn
           (\, crx \, (\, q \,, \ p \, [m \, \% \,\, n] \,\,, \ p \, [\, (m \,+\, 1) \,\, \% \,\, n \,] \,) \, = \!\!\!\!\! - s \,) \ l \, = m;
      \begin{array}{ll} \textbf{else} & r = m; \end{array}
   update_tangent(q, r % n, a, b);
bool contain (P p) {
   if (p.x < d[0].x | | p.x > d.back().x) return 0;
   auto it
        = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
   if (it->x = p.x) {
   if (it->y > p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;
   it = lower_bound
         (u.begin(), u.end(), P(p.x, 1e12), greater<P>());
   if (it->x = p.x) {
   if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
   a = b = 0;
   int i
        = lower_bound(d.begin(), d.end(), p) - d.begin();
  \begin{array}{l} bs(\,0\,,\ i\,,\ p\,,\ a\,,\ b\,)\,;\\ bs(\,i\,,\ d\,.\,size\,(\,)\,,\ p\,,\ a\,,\ b\,)\,; \end{array}
   i = lower\_bound(
        u.begin(), u.end(), p, greater<P>()) - u.begin();
   bs((int
         d.size() - 1, (int)d.size() - 1 + i, p, a, b);
   bs((int)d.size()
            1 + i, (int)d.size() - 1 + u.size(), p, a, b);
};
```

8.15 3D Convex Hull [29e4a9]

```
absvol(const P a, const P b, const P c, const P d) {
  return abs(((b-a)^(c-a))*(d-a))/6;
struct convex3D {
  static const int maxn=1010;
  struct T{
    int a,b,c;
    bool res;
    T()\{\}
    T(int a, int
          b, int c, bool res=1):a(a),b(b),c(c),res(res){}
  int n,m;
 P p [maxn];
 T f [maxn*8];
  int id [maxn] [maxn];
  bool on (T &t, P &q) {
    return ((
         p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
  int g=id[a][b];
    if (f[g].res){
       if(on(f[g],p[q]))dfs(q,g);
         id[q][b]=id[a][q]=id[b][a]=m;
         f[m++]=T(b,a,q,1);
    }
  void dfs(int p,int i){
    f[i].res=0;
    meow(p, f[i].b, f[i].a);

meow(p, f[i].c, f[i].b);
    meow(p, f[i].a, f[i].c);
  void operator()(){
    if (n<4)return;
    if ([&](){
         for (int i=1;i< n;++i) if (abs
             (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
         return 1;
}() || [&](){
         for (int
                  i=2; i < n; ++i) if (abs((p[0]-p[i])
              ^(p[1]-p[i])>eps)return swap(p[2],p[i]),0;
         return 1
         }() || [&](){
```

```
for (int i
                                                       = 3; i < n; ++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                                                       *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
                                    }())return;
                   for (int i=0; i<4;++i){
                          T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
                            if(on(t,p[i]))swap(t.b,t.c);
                            id [t.a][t.b]=id [t.b][t.c]=id [t.c][t.a]=m;
                            f[m++]=t;
                   for (int i=4; i< n; ++i) for
                                     (int j=0; j < m++j) if (f[j].res & on(f[j],p[i])) {
                            dfs(i,j);
                            break;
                   int mm=m; m=0;
                   for (int i=0; i \le mm + +i) if (f[i].res) f [m++]=f[i];
          bool same(int i, int j)
                  return !(absvol(p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].c]),p[f[i].c],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[
                                     (p[f[i].a], p[f[i].b], p[f[i].c], p[f[j].c]) > eps);
          int faces(){
                   int r=0;
                   for (int i=0; i \triangleleft m; ++i){
                            int iden=1;
                            for (int j=0; j< i; ++j) if (same(i,j)) iden=0;
                           r += iden;
                   return r;
} tb;
```

8.16 Minimum Enclosing Circle [fc0e72]

```
pt center (const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^
                    p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {
        (norm2(cent - p[i]) \le r) continue;
     cent = p[i];
     r = 0.0;
     for (int j = 0; j < i; ++j) {
       if (norm2(cent - p[j]) \le r) continue;
       cent = (p[i] + p[j])'/r = norm2(p[j] - cent);
       for (int k = 0; k < j; ++k) {
         if (norm2(cent - p[k]) <= r) continue;</pre>
         cent = center(p[i], p[j], p[k]);
         r = norm2(p[k] - cent);
    }
  }
  return circle(cent, sqrt(r));
```

8.17Closest Pair [f6de57]

```
double closest_pair(int l, int r) {
  // p should be sorted
        increasingly according to the x-coordinates.
     (1 = r) return 1e9;
  if (r - l = 1) return dist(p[l], p[r]);
  int m = (l + r) >> 1;
  double d =
       min(closest_pair(l, m), closest_pair(m + 1, r));
  vector<int> vec;
  \quad \text{for (int } i = m; \ i >= 1 \ \&\& \\
       fabs(p[m].x - p[i].x) < d; --i) \ vec.push\_back(i);
      (int i = m + 1; i \le r \&\&
       fabs(p[m].x - p[i].x) < d; ++i) vec.push_back(i);
  sort(vec.begin(), vec.end()
        [\&](int \ a, \ int \ b) \ \{ \ return \ p[a].y < p[b].y; \ \});
  for (int i = 0; i < vec.size(); ++i) {
```

```
for (int j = i + 1; j < vec.size()
    && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {
    d = min(d, dist(p[vec[i]], p[vec[j]]));
    }
}
return d;
}</pre>
```

9 Else

9.1 Cyclic Ternary Search* [9017cc]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

9.2 Mo's Algorithm (With modification) $_{[f05c5b]}$

```
Mo's Algorithm With modification
 Block: N^{2/3}, Complexity: N^{5/3}
 struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
    L(1), R(r), LBid(1 / blk), RBid(r / blk), T(t) {}
bool operator < (const Query &q) const {
        \quad \text{if } (LBid \mathrel{!=} q.LBid) \enspace \textbf{return} \enspace LBid < q.LBid;
        if (RBid != q.RBid) return RBid < q.RBid;
        return T < b.T;
 };
 void solve(vector<Query> query) {
    sort(ALL(query));
    int L=0, R=0, T=-1;
    \begin{array}{lll} & for \ (auto \ q \ : \ query) \ \{ \\ & while \ (T < q.T) \ addTime(L, \ R, \ +\!\!+\!\!T); \ /\!/ \ TODO \end{array}
        while (T > q.T) subTime(L, R, T--); // TODO
        while (R < q.R) add(arr[++R]); // TODO while (L > q.R) add(arr[--L]); // TODO while (R > q.R) sub(arr[R--]); // TODO while (L < q.L) sub(arr[L++]); // TODO
        // answer query
}
```

9.3 Mo's Algorithm On Tree [8331c2]

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord [in[u]] = ord [out[u]] = u
4) bitset MAXN> inset
struct Query
  int L, R, LBid, lca;
  Query(int u, int v) {
     int c = LCA(u, v);
     if (c = u \mid \mid c = v)
       q.\,lc\,a\,=\,-1\,,\,\,q.\,L\,=\,out\,[\,c\,\,\widehat{}\,\,u\,\,\widehat{}\,\,v\,]\,\,,\,\,q.\,R\,=\,out\,[\,c\,\,]\,;
     else if (out[u] < in[v])
       q.lca = c, q.L = out[u], q.R = in[v];
     else
       q.lca = c, q.L = out[v], q.R = in[u];
    q.Lid = q.L / blk;
  bool operator < (const Query &q) const {
    if (LBid != q.LBid) return LBid < q.LBid;
     return R < q.R;
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
```

```
inset[x] = ~inset[x];
}
void solve(vector<Query> query) {
    sort(ALL(query));
    int L = 0, R = 0;
    for (auto q : query) {
        while (R < q.R) flip(ord[++R]);
        while (L > q.L) flip(ord[--L]);
        while (R > q.R) flip(ord[R--]);
        while (L < q.L) flip(ord[L++]);
        if (~q.lca) add(arr[q.lca]);
        // answer query
        if (~q.lca) sub(arr[q.lca]);
    }
}</pre>
```

9.4 Additional Mo's Algorithm Trick

- Mo's Algorithm With Addition Only
 - Sort queryssame as the normal Mo's algorithm.
 - For each query [l,r]:
 - If l/blk = r/blk, brute-force.
 - If $\dot{l}/blk \neq curL/blk$, initialize $curL := (l/blk + 1) \cdot blk$, curR := curL 1
 - If r > curR, increase curR
 - decrease curL to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
 - Part1: Answer all f([1,r],r+1) first.
 - Part2: Store $curR \to R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
 - Note: You must do the above symmetrically for the left boundaries.

9.5 Hilbert Curve [1274a3]

```
1ll hilbert(int n, int x, int y) {
    ll res = 0;
    for (int s = n / 2; s; s >>= 1) {
        int rx = (x & s) > 0;
        int ry = (y & s) > 0;
        res += s * 1ll * s * ((3 * rx) ^ ry);
        if (ry == 0) {
            if (rx == 1) x = s - 1 - x, y = s - 1 - y;
            swap(x, y);
        }
    }
    return res;
} // n = 2^k
```

9.6 DynamicConvexTrick* [673ffd]

```
only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
      <(const Line &rhs) const { return a < rhs.a; }</pre>
  bool operator <(ll x) const { return p < x; }
struct DynamicHull : multiset<Line, less >> {
  static const ll kInf = 1e18;
  if (y = end()) \{ x->p = kInf; return 0; \}
    if (x
        ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
    else x->p = Div(y->b - x->b, x->a - y->a);
    return x->p>=y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
    () && isect(--x, y)) isect(x, y = erase(y));
while ((y = x) != begin
        () && (--x)-p >= y-p isect(x, erase(y));
  ll query(ll x) {
    auto l = *lower_bound(x);
    return l.a * x + l.b;
};
```

9.7 All LCS* [78a378]

9.8 AdaptiveSimpson* [4074b3]

```
template<typename Func, typename d = double>
struct Simpson {
  using pdd = pair < d, d>;
 Func f;
 d eval(pdd l, pdd r, d fm, d eps) {
   pdd m((1.X + r.X) / 2, fm);
    d = mix(1, r, fm) . second;
   auto [flm, sl] = mix(l, m);
     auto [fmr, sr] = mix(m, r); 
    d \ delta = sl + sr - s;
    if (abs(delta
        ) <= 15 * eps) return sl + sr + delta / 15;
    return eval(l, m, flm, eps / 2) +
      eval(m, r, fmr, eps / 2);
 d eval(d l, d r, d eps) {
    return eval
        (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
 d eval2(d l, d r, d eps, int k = 997) {
  d h = (r - l) / k, s = 0;
    for (int i = 0; i < k; ++i, l + s; s + eval(1, 1 + h, eps / k);
                                l += h
    return s;
};
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
```

9.9 Simulated Annealing [de78c6]

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.99995;
}
```

9.10 Tree Hash* [34aae5]

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
            sum += shift(dfs(i, u));
    return sum;
}</pre>
```

9.11 Binary Search On Fraction [765c5a]

```
struct Q {
    ll p, q;
    Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
};
bool pred(Q);
```

```
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N) {
Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
        if (Q mid = hi.go(lo, len + step);
            mid.p > N || mid.q > N || dir ^ pred(mid))
            t++;
        else len += step;
        swap(lo, hi = hi.go(lo, len));
        (dir ? L : H) = !!len;
    }
    return dir ? hi : lo;
}
```

9.12 Bitset LCS [330ab1]

```
 \begin{array}{l} \mbox{cin} >> n >> m; \\ \mbox{for (int } i = 1, \ x; \ i <= n; \ +\!\!+\!\!i) \\ \mbox{cin} >> x, \ p[x].set(i); \\ \mbox{for (int } i = 1, \ x; \ i <= m; \ i\!\!+\!\!+\!\!) \left\{ \\ \mbox{cin} >> x, \ (g = f) \ |= p[x]; \\ \mbox{f.shiftLeftByOne(), f.set(0);} \\ \mbox{((f = g - f) } \widehat{\ }= g) \ \&= g; \\ \mbox{cout} << f.count() << '\n'; \\ \end{array}
```

9.13 N Queens Problem [dlfccc]

```
void solve
  (vector<int> &ret, int n) { // no sol when n=2,3
  if (n % 6 == 2) {
    for (int i = 2; i <= n; i += 2) ret.pb(i);
    ret.pb(3); ret.pb(1);
    for (int i = 7; i <= n; i += 2) ret.pb(i);
    ret.pb(5);
  } else if (n % 6 == 3) {
    for (int i = 4; i <= n; i += 2) ret.pb(i);
    ret.pb(2);
    for (int i = 5; i <= n; i += 2) ret.pb(i);
    ret.pb(1); ret.pb(3);
  } else {
    for (int i = 2; i <= n; i += 2) ret.pb(i);
    for (int i = 1; i <= n; i += 2) ret.pb(i);
  }
}</pre>
```

10 Python

10.1 Misc