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region and then type : Hash to hash your selection."

"Select

```
1.2 readchar [a419b9]
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
  static char *p = buf, *end = buf;
  if (p = end) end = buf +
       fread_unlocked(buf, 1, bufsize, stdin), p = buf;
1.3 BigIntIO [d9afcb]
  _int128 read() {
       int128 \ x = 0, f = 1;
    char ch = getchar();
    while (ch < '0', | | ch > '9') {
    if (ch == '-') f = -1;
         ch = getchar();
    while (ch >= '0' && ch <= '9') {
    x = x * 10 + ch - '0';
         ch = getchar();
    return x * f;
void print(__int128 x) {
    if (x < 0) {
putchar('-');
        x = -x:
    if (x > 9) print(x / 10);
    putchar(x % 10 + '0'');
bool cmp(\underline{\phantom{a}}int128 x, \underline{\phantom{a}}int128 y) { return x > y; }
1.4 Black Magic [0d8b5f]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
     tree<int, null_type, std::less<int>, rb_tree_tag
      tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int , int> umap;
typedef priority_queue<int> heap;
int main() {
  // random
  mt19937 rng(chrono::
       steady_clock::now().time_since_epoch().count());
      get_rand(int l, int r){ return
        uniform_int_distribution<int>(l, r)(rng); }
  shuffle(v.begin(), v.end(), rng);
  // rb tree
  tree set s
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
      (0) = 22; assert(*s.find_by_order(1) = 71);
  assert (s.order_of_key
      (22) = 0; assert (s.order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by_order
       (0) = 71; assert (s.order_of_key(71) = 0);
   // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope < char > r[2];
  \mathbf{r}[1] = \mathbf{r}[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0;
```

1.5 Pragma Optimization [eac636]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
 \_builtin\_ia32\_ldmxcsr(\_builtin\_ia32\_stmxcsr()|0x8040)
```

Bitset [cb5d05]

```
#include < bits / stdc++.h>
using namespace std;
int main () {
     bitset <4> bit;
     bit.all(); // all bit is true, ret tru;
bit.any(); // any bit is true, ret true
     bit.none(); // all bit is false, ret true
     bit.count();
     bit.to_string('0', '1');//with parmeter
     bit.reset(); // set all to true
bit.set(); // set all to false
     std::bitset <%> b3{0}, b4{42};
     std::hash<std::bitset<8>> hash_fn8;
     hash_fn8(b3); hash_fn8(b4);
}
```

$\mathbf{2}$ Graph

2.1 BCC Vertex* [740acb]

```
struct BCC { // 0-base
  int n, dft, nbcc;
   vector<int> low, dfn, bln, stk, is_ap, cir;
   vector<vector<int>>> G, bcc, nG;
   void make_bcc(int_u) {
      bcc.emplace_back(1, u);
     for (; stk.back() != u; stk.pop_back())
  bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
stk.pop_back(), bln[u] = nbcc++;
   void dfs(int u, int f) {
      int child = 0;
      \begin{array}{ll} low [u] = dfn [u] = ++dft \,, \; stk.pb(u) \,; \\ for \; (int \; v \; : \; G[u]) \end{array}
         if (!dfn[v]) {
            dfs(v, u), ++child; low[u] = min(low[u], low[v]);
            \begin{array}{l} \mbox{if } (d f n [u] <= low[v]) \ \{ \\ \mbox{is\_ap}[u] = 1, \ b l n [u] = n b c c \,; \end{array}
               make\_bcc(v), bcc.back().pb(u);
         \} \ \ else \ \ if \ \ (dfn[v] < dfn[u] \ \&\& \ v \ != \ f)
      \begin{array}{l} low[u] = min(low[u], dfn[v]); \\ if (f = -1 \&\& child < 2) is_ap[u] = 0; \end{array}
      if (f = -1 \&\& child = 0) make\_bcc(u);
  BCC(int _n): n(_n), dft(),
   G[u].pb(v), G[v].pb(u);
  if (!dfn[i]) dfs(i, -1);
   void block_cut_tree() {
      cir.resize(nbcc);
      for (int i = 0; i < n; ++i)
         if (is_ap[i])
            bln[i] = nbcc++;
      cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
         for (int j : bcc[i])
            if (is_ap[j])
  \begin{array}{c} n\ddot{G}\,[\,\,\bar{i}\,\,].\,pb(\,bln\,[\,j\,\,])\,\,,\,\,nG[\,bln\,[\,j\,\,]]\,.\,pb(\,i\,)\,;\\ \}\,\,//\,\,up\,\,to\,\,2\,\,*\,\,n\,\,-\,\,2\,\,nodes\,!!\,\,\,bln\,[\,i\,\,]\,\,for\,\,id \end{array}
```

2.2 Bridge* [4da29a]

```
struct ECC { // 0-base
  int n, dft, ecnt, necc;
vector<int> low, dfn, bln, is_bridge, stk;
  vector<vector<pii>>> G;
  void dfs(int u, int f) {
    dfn[u] = low[u] = ++dft, stk.pb(u);
    for (auto [v, e] : G[u])
```

```
if (!dfn[v])
      dfs(v, e), low[u] = min(low[u], low[v]);
else if (e != f)
    low[u] = min(low[u], dfn[v]);
if (low[u] == dfn[u]) {
         if (f!= -1) is_bridge[f] = 1;
for (; stk.back() != u; stk.pop_back())
           bln[stk.back()] = necc;
         bln[u] = necc++, stk.pop\_back();
     }
   ÉCC(int _n): n(_n), dft()
   , ecnt(), necc(), low(n), dfn(n), bln(n), G(n) {} void add_edge(int u, int v) {}
     G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
   void solve() {
      is_bridge.resize(ecnt);
      for (int i = 0; i < n; ++i)
         if (!dfn[i]) dfs(i, -1);
}; // ecc_id(i): bln[i]
2.3 SCC* [4057dc]
struct SCC { // 0-base
   int n, dft, nscc;
   vector <int > low, dfn, bln, instack, stk;
   vector<vector<int>>> G;
   void dfs(int u)
      low[u] = dfn[u] = ++dft;
      instack[u] = 1, stk.pb(u);
      for (int v : G[u])
         if (!dfn[v])
         d\hat{f}s(v), low[u] = min(low[u], low[v]);

else if (instack[v] && dfn[v] < dfn[u])
           low[u] = min(low[u], dfn[v]);
      if (low u = dfn[u]) {
  for (; stk.back() != u; stk.pop_back())
           bln[stk
         \label{eq:back()} \begin{array}{l} \left. \operatorname{back}() \right] = \operatorname{nscc}, \ \operatorname{instack}\left[\operatorname{stk.back}()\right] = 0; \\ \operatorname{instack}\left[u\right] = 0, \ \operatorname{bln}\left[u\right] = \operatorname{nscc++}, \ \operatorname{stk.pop\_back}(); \end{array}
  void add_edge(int u, int v) {
     G[u].pb(v);
```

}; // scc_id(i): bln[i] 2.4 2SAT* [f5630a]

for (int i = 0; i < n; ++i)

if (!dfn[i]) dfs(i);

void solve() {

```
struct SAT { // 0-base
   int n;
   vector<bool> istrue;
   SCC scc;
  \begin{array}{lll} SAT(\ int \ \_n): \ n(\_n) \ , \ istrue(n+n) \ , \ scc(n+n) \ \{\} \\ int \ rv(\ int \ a) \ \{ \end{array}
     return a > = n? a - n : a + n;
   void add_clause(int a, int b) {
     scc.add\_edge(rv(a)\,,\ b)\,,\ scc.add\_edge(rv(b)\,,\ a)\,;
   bool solve()
     scc.solve();
      for (int i = 0; i < n; ++i) {
        if (scc.bln[i] = scc.bln[i + n]) return false;
istrue[i] = scc.bln[i] < scc.bln[i + n];
        istrue[i + n] = !istrue[i];
     return true;
};
```

2.5 MinimumMeanCycle* [3e5d2b]

```
ll road[N][N]; // input here
struct MinimumMeanCycle {
  11\ dp\,[N\,+\,\,5\,]\,[N]\ ,\ n\,;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
     for (int i = 2; i \le L; ++i)
```

```
for (int k = 0; k < n; ++k)
    for (int j = 0; j < n; ++j)
        dp[i][j] =
        min(dp[i - 1][k] + road[k][j], dp[i][j]);

for (int i = 0; i < n; ++i) {
    if (dp[L][i] >= INF) continue;
    ll ta = 0, tb = 1;
    for (int j = 1; j < n; ++j)
        if (dp[j][i] < INF &&
        ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
        ta = dp[L][i] - dp[j][i], tb = L - j;
    if (ta == 0) continue;
    if (a == -1 | | a * tb > ta * b) a = ta, b = tb;
}

if (a != -1) {
    ll g = __gcd(a, b);
    return pll(a / g, b / g);
}

return pll(-1LL, -1LL);
}

void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
}
};</pre>
```

2.6 Virtual Tree* [1b641b]

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
  if (top = -1) return st[++top] = u, void();
  int p = LCA(st[top], u);
  if (p = st[top]) return st[++top] = u, void();
  while (top >= 1 \&\& dep[st[top - 1]] >= dep[p])
  vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
  \quad \text{for (int`i : vG[u]) reset(i);} \\
 vG[u].clear();
void solve(vector<int> &v) {
  top = -1
  sort (ALL(v),
    [\,\&\,](\,int\ a\,,\ int\ b\,)\ \{\ return\ dfn\,[\,a\,]\,<\,dfn\,[\,b\,]\,;\ \})\,;
  for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do something
  reset(v[0]);
```

2.7 Maximum Clique Dyn* [d50aa9]

```
struct MaxClique { // fast when N \le 100
  bitset \triangleleft \triangleright G[N], cs[N];
int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = \underline{n};
    for (int i = 0; i < n; ++i) G[i].reset();
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<№ mask) {
    if (1 < 4) {
       for (int i: r) d[i] = (G[i] \& mask).count();
      sort (ALL(r)
            (x, [x](int x, int y) \{ return d[x] > d[y]; \});
    vector < int > c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
for (int p : r) {
      int k = 1:
       while ((cs[k] \& G[p]).any()) ++k;
       if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
       if(k < lft) r[tp++] = p;
    for (int k = lft; k \ll rgt; ++k)
```

```
for (int p = cs[k]._Find_first
         (); p < N; p = cs[k]._Find_next(p))
r[tp] = p, c[tp] = k, ++tp;
     dfs(r, c, l + 1, mask);
  void dfs(vector<
       int > &r, vector < int > &c, int l, bitset < N > mask) {
     while (!r.empty()) {
       int p = r.back();
       r.pop\_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;</pre>
       {\rm cur}\,[\, q{+}{+}]\,=\,p\,;
       vector<int> nr;
       for (int i : r) if (G[p][i]) nr.pb(i);
       if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
       else if (q > ans) ans = q, copy_n(cur, q, sol);
       c.pop_back(), --q;
  int solve() {
    vector < int > r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset < N > (string(n, '1')));
    return ans;
};
```

2.8 Minimum Steiner Tree* [62d6fb]

```
struct SteinerTree { // 0-base
  \begin{array}{l} \text{int } n, \ dst\left[N\right]\left[N\right], \ dp\left[1 << T\right]\left[N\right], \ tdst\left[N\right]; \\ \text{int } vcst\left[N\right]; \ // \ the \ cost \ of \ vertexs \end{array}
  void init(int _n) {
     n = \underline{n};
     for (int i = 0; i < n; ++i) {
fill_n(dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
    x = \min(x, val);
  void add_edge(int ui, int vi, int wi) {
     chmin (dst [ui][vi], wi);
  void shortest_path() {
     for (int k = 0; k < n; ++k)
       for (int i = 0; i < n; ++i)
          for (int j = 0; j < n; ++j)

chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1;
     for (int i = 0; i \le full; ++i)
       fill_n (dp[i], n, INF);
     copy_n(vcst, n, dp[0]);
for (int msk = 1; msk <= full; ++msk) {</pre>
       if (!(msk & (msk - 1))) {
          dp [msk
                  [i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; ++i)
          for (int sub = (
               msk - 1) \& msk; sub; sub = (sub - 1) \& msk)
            chmin (dp [msk] [i]
                 dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
        for (int i = 0;
                           i < n; ++i) {
          tdst[i] = INF;
          for (int j = 0; j < n; ++j)
            chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
; // O(V 3^T + V^2 2^T)
```

2.9 Dominator Tree* [2b8b32]

```
void init(int _n) {
  n = _n;
for (int i = 1; i <= n; ++i)
     G[i].clear(), rG[i].clear();
 \begin{array}{c} \textbf{void} \ \ \textbf{add\_edge(int} \ \ \textbf{u}, \ \ \textbf{int} \ \ \textbf{v}) \ \ \{\\ G[\textbf{u}].\, pb(\textbf{v}) \,, \ rG[\textbf{v}].\, pb(\textbf{u}) \,; \end{array} 
void dfs(int u) {
  id [dfn[u] = ++Time] = u;
for (auto v : G[u])
     if(!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
int find(int y, int x) {
  if (y \le x) return y;
   int tmp = find(pa[y], x);
   if \ (semi[best[y]] > semi[best[pa[y]]]) \\
     best[y] = best[pa[y]];
  return pa[y] = tmp;
void tarjan(int root) {
  Time = 0;
   for (int i = 1; i \le n; ++i) {
     dfn[i] = idom[i] = 0;
     tree[i].clear();
     best[i] = semi[i] = i;
   dfs(root);
   for (int i = Time; i > 1; --i) {
     int u = id[i];
     for (auto v : rG[u])
if (v = dfn[v]) {
          find(v, i);
semi[i] = min(semi[i], semi[best[v]]);
     tree[semi[i]].pb(i);
for (auto v : tree[pa[i]]) {
        find(v, pa[i]);
idom[v] =
           semi[best[v]] = pa[i] ? pa[i] : best[v];
     tree [pa[i]]. clear();
   for (int i = 2; i \leftarrow Time; ++i)
     if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
     tree[id[idom[i]]].pb(id[i]);
```

2.10 Four Cycle [584a52]

```
int main() {
  cin.tie(nullptr)->sync_with_stdio(false);
  cin >> n >> m;
  for (int i = 1; i \le m; i++) {
    \quad \quad \text{int} \ u\,,\ v\,;
    cin >> u >> v;
    E[\,u\,]\,.\,push\_back\,(\,v\,)\,;
    E[v].push_back(u);
    deg[u]++, deg[v]++;
  for (int u = 1; u \le n; u++)
    for (int v : E[u])

if (\deg[u] > \deg[v] \mid | (\deg[u] \cdot push\_back(v);

[u] = \deg[v] & u > v) E1[u].push_back(v);
  for (int a = 1; a \le n; a++) {
    for (int b : E1[a])
       for (int c : E[b]) {
         total += cnt[c]++;
    for (int b : E1[a])
      for (int c : E[b]) cnt[c] = 0;
  cout << total << '\n';
  return 0;
```

2.11 Minimum Clique Cover* [879472]

```
fill_n(E, n, 0), fill_n(co, 1 << n, 0);
}
void add_edge(int u, int v) {
    E[u] |= 1 << v, E[v] |= 1 << u;
}
int solve() {
    for (int i = 0; i < n; ++i)
        co[1 << i] = E[i] | (1 << i);
    co[0] = (1 << n) - 1;
    dp[0] = (n & 1) * 2 - 1;
    for (int i = 1; i < (1 << n); ++i) {
        int t = i & -i;
        dp[i] = -dp[i ^ t];
        co[i] = co[i ^ t] & co[t];
}
for (int i = 0; i < (1 << n); ++i)
        co[i] = (co[i] & i) == i;
    fwt(co, 1 << n, 1);
    for (int ans = 1; ans < n; ++ans) {
        int sum = 0; // probabilistic
        for (int i = 0; i < (1 << n); ++i)
            sum += (dp[i] *= co[i]);
        if (sum) return ans;
}
return n;
}
</pre>
```

2.12 NumberofMaximalClique* [11fa26]

```
struct BronKerbosch { // 1-base
   \begin{array}{l} int \ n, \ a[N] \ , \ g[N] \ [N] \ ; \\ int \ S, \ all \ [N] \ [N] \ , \ some \ [N] \ [N] \ , \ none \ [N] \ [N] \ ; \end{array}
   void init(int _n) {
      \begin{array}{lll} n = \underline{\ \ } n; \\ \mbox{for (int } i = 1; \ i <= n; \ +\!\!\!+\!\! i) \end{array}
          for (int j = 1; j \le n; ++j) g[i][j] = 0;
   void add_edge(int u, int v) {
      g[u][v] = g[v][u] = 1;
    void dfs(int d, int an, int sn, int nn) {
      if (S > 1000) return; // pruning
if (sn == 0 && nn == 0) ++S;
      int u = some[d][0];
       for (int i = 0; i < sn; ++i) {
          int v = some[d][i];
          if (g[u][v]) continue;
int tsn = 0, tnn = 0;
          copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
          for (int j = 0; j < sn; ++j)
  if (g[v][some[d][j]])</pre>
                some[d + 1][tsn++] = some[d][j];
          for (int j = 0; j < nn; ++j)
if (g[v][none[d][j]])
                none[d + 1][tnn++] = none[d][j];
          dfs(d + 1, an + 1, tsn, tnn);

some[d][i] = 0, none[d][nn++] = v;
      }
   int solve() {
  iota(some[0], some[0] + n, 1);
  S = 0, dfs(0, 0, n, 0);
       return S;
```

3 Data Structure

int bit [N + 1]; //N = 2 ^ k int query_kth(int k) {

3.1 Discrete Trick

```
int res = 0:
    for (int i = N >> 1; i >= 1; i >>= 1)
        if (bit[res + i] < k)
            k -= bit[res += i];
    return res + 1;
}
```

3.3 Interval Container* [c54d29]

```
/* Add and
      remove intervals from a set of disjoint intervals.
   Will merge the added interval with
        any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>>::
     iterator addInterval(set<pii> is, int L, int R) {
  \begin{array}{ll} \mbox{if } (L \Longrightarrow R) \mbox{ return is.end();} \\ \mbox{auto it} = \mbox{is.lower\_bound(\{L,\,R\}), before} = \mbox{it}; \\ \end{array}
  while (it != is.end() && it->X <= R) {
    R = \max(R, \ it \text{->}Y);
     before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y>= L) {
    L = \min(L, it ->X);
    R = \max(R, it ->Y);
     is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii>% is, int L, int R) {
  if (L = R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it ->Y;
  if (it->X == L) is.erase(it);
   \begin{array}{ll} \textbf{else} & (\texttt{int\&}) \texttt{it} \text{-} \texttt{>} Y = L; \\ \end{array} \\
  if (R != r2) is .emplace(R, r2);
```

3.4 Leftist Tree [e91538]

```
struct node {
  ll v, data, sz, sum;
node *1, *r;
  node(ll k)
     : v(0), data(k), sz(1), l(0), r(0), sum(k) \{ \}
f,
ll sz(node *p) { return p ? p->sz : 0; }
ll V(node *p) { return p ? p->v : -1; }
ll sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
}
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
   delete tmp;
```

3.5 Heavy light Decomposition* [b004ae]

```
\begin{array}{c} \textbf{struct Heavy\_light\_Decomposition} \ \{ \ // \ \ 1\text{-}\textbf{base} \\ \textbf{int} \ \ n, \ \ ulink \ [N] \ , \ \ deep \ [N] \ , \ mxson \ [N] \ , \ \ w[N] \ , \ pa \ [N] \ ; \end{array}
   int t, pl[N], data[N], val[N]; // val: vertex data
   vector < int > G[N];
   void init(int _n) {
     \begin{array}{l} n = \_n; \\ \text{for (int } i = 1; \ i <= n; \ +\!\!+\!\! i) \end{array}
         G[i]. clear(), mxson[i] = 0;
   void add_edge(int a, int b) {
     G[a].pb(b), G[b].pb(a);
   void dfs(int u, int f, int d) {
     w[u] = 1, pa[u] = f, deep[u] = d++;
      for (int &i : G[u])
         if (i != f)
             dfs(i, u, d), w[u] += w[i];
             i\,f\ (w[\,mxson\,[\,u\,]\,]\,<\,w[\,i\,]\,)\ mxson\,[\,u\,]\,=\,i\,;
   void cut(int u, int link) {
```

```
if (!mxson[u]) return;
    cut(mxson[u], link);
for (int i : G[u])
       if (i != pa[u] && i != mxson[u])
         cut(i, i);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
  int ta = ulink[a], tb = ulink[b], res = 0;
     while (ta != tb) {
       if (deep
            [ta] > deep[tb] swap(ta, tb), swap(a, b);
        / \text{ query(pl[tb], pl[b])}
       tb = ulink[b = pa[tb]];
     if (pl[a] > pl[b]) swap(a, b);
    // query(pl[a], pl[b])
};
```

3.6 Centroid Decomposition* [5a24da]

```
struct Cent_Dec { // 1-base
  vector<pll> G[N];
pll info[N]; // store info. of itself
pll upinfo[N]; // store info. of climbing up
int n, pa[N], layer[N], sz[N], done[N];
ll dis[_lg(N) + 1][N];
   void init (int _n) {
     n = \underline{n}, layer[0] = -1;
fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
      for (int i = 1; i \le n; ++i) G[i]. clear();
   void add_edge(int a, int b, int w) {
     G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
   void get_cent(
      int u, int f, int &mx, int &c, int num) {
      int mxsz = 0;
      sz[u] = 1;
      for (pll e : G[u])
          \begin{array}{lll} & \text{if } (!\, done\, [\, e . \dot{X}] & \&\& \ e . X \ != \ f) \ \{ \\ & \text{get\_cent}\, (\, e . X, \ u, \ mx, \ c, \ num) \, ; \\ & \text{sz}\, [\, u] \ += \ sz\, [\, e . X] \, , \ mxsz \ = \ max(mxsz \, , \ sz\, [\, e . X]) \, ; \\ \end{array} 
       i \, f \, \left( mx > \, max(\, mxsz \, , \, \, num \, - \, \, sz \, [\, u \, ] \, \right) \, ) 
        mx = max(mxsz, num - sz[u]), c = u;
   void dfs(int u, int f, ll d, int org) {
      // if required, add self info or climbing info dis[layer[org]][u] = d;
      for (pll e : G[u])
if (!done[e.X] && e.X!= f)
            dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
      int mx = 1e9, c = 0, lc;
      get_cent(u, f, mx, c, num);
     done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
for (pll e : G[c])
if (!done[e.X]) {
            if (sz[e.X] > sz[c])
            lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
            upinfo\,[\,lc\,] \,=\, pll\,(\,)\;,\;\; dfs\,(\,e\,.X,\;\; c\,,\;\; e\,.Y,\;\; c\,)\;;
     return done [c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
      for (int a = u, ly = layer[a]; a;
             a = pa[a], --ly)
         info[a].X += dis[ly][u], ++info[a].Y;
         if (pa[a])
            upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
  11 query(int u) {
     11 \text{ rt} = 0;
      for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) \{

rt += info[a].X + info[a].Y * dis[ly][u];
         if (pa[a])
```

 $upinfo\left[\left.a\right].X+\,upinfo\left[\left.a\right].Y\,\,*\,\,dis\left[\left.ly\,\,-\,\,1\right]\left[u\right];$

for (Splay *q = x;; q = q->f) {

```
splayVec.pb(q);
      return rt:
                                                                                        if (q->isr()) break;
};
                                                                                     reverse (ALL(splayVec));
      LiChaoST* [4a4bee]
3.7
                                                                                     for (auto it : splayVec) it->push();
                                                                                     while (!x->isr()) {
   if (x->f->isr()) rotate(x);
struct L {
    ll m, k, id;
                                                                                        else if (x->dir() = x->f->dir())
  L() : id(-1) \{ \}
                                                                                           rotate(x->f), rotate(x);
  \begin{array}{l} L(\hat{1}l\ a,\ \hat{l}l\ b,\ \hat{l}l\ c):m(a),\ k(b),\ id(c)\ \{\}\\ ll\ at(ll\ x)\ \{\ return\ m\ *\ x+k;\ \} \end{array}
                                                                                        else rotate(x), rotate(x);
class LiChao { // maintain max
                                                                                  Splay* access(Splay *x) {
private:
                                                                                     Splay *q = nil;
for (; x != nil; x = x->f)
   int n; vector<L> nodes;
  void insert(int 1, int r, int rt, L ln) {
  int m = (1 + r) >> 1;
  if (nodes[rt].id == -1)
                                                                                        \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
                                                                                     return q;
        return nodes[rt] = ln, void();
                                                                                   void root_path(Splay *x) { access(x), splay(x); }
      bool atLeft = nodes[rt].at(1) < ln.at(1);
                                                                                  void chroot(Splay *x){
      if (nodes[rt].at(m) < ln.at(m))
                                                                                     root_path(x), x->give_tag(1);
      atLeft = 1, swap(nodes[rt], ln);
if (r - l == 1) return;
                                                                                     x \rightarrow push(), x \rightarrow pull();
      if (atLeft) insert(l, m, rt << 1, ln);</pre>
                                                                                  void split (Splay *x, Splay *y) {
      else insert (m, r, rt \ll 1 \mid 1, ln);
                                                                                     chroot(x), root_path(y);
   11 query(int l, int r, int rt, ll x) {
                                                                                  void link (Splay *x, Splay *y) {
      int m = (l + r) \gg 1; ll ret = -INF;
                                                                                     root\_path(x), chroot(y);
      if (nodes[rt].id != -1) ret = nodes[rt].at(x);
                                                                                     x-\operatorname{setCh}(y, 1);
          (r - l = 1) return ret;
      if (x
                                                                                  void cut(Splay *x, Splay *y) {
            < m) return max(ret, query(l, m, rt << 1, x));
                                                                                     split(x, y);
       return max(ret, query(m, r, rt << 1 | 1, x)); 
                                                                                     if (y->size != 5) return;
                                                                                     y->push();
public:
                                                                                     y->ch[0] = y->ch[0]->f = nil;
  LiChao(int n_) : n(n_), nodes(n * 4) {}
   void insert(L ln) { insert(0, n, 1, ln); }
                                                                                  Splay* get_root(Splay *x)
   ll query(ll x) { return query(0, n, 1, x); }
                                                                                     for (\text{root\_path}(x); x->\text{ch}[0] != \text{nil}; x = x->\text{ch}[0])
                                                                                       x \rightarrow push();
                                                                                     splay(x);
3.8 Link cut tree* [a35b5d]
                                                                                     return x;
struct Splay { // xor-sum
                                                                                  bool conn(Splay *x, Splay *y) {
   static Splay nil;
Splay *ch[2], *f;
                                                                                     return get_root(x) == get_root(y);
   int val, sum, rev, size;
                                                                                  Splay* lca(Splay *x, Splay *y) {
   Splay (int
                                                                                     access(x), root_path(y);
        _{\text{val}} = 0 : val(_{\text{val}}), sum(_{\text{val}}), rev(0), size(1)
                                                                                     if (y->f = nil) return y;
   \{ f = ch[0] = ch[1] = &nil; \}
                                                                                     return y -> f;
   bool isr()
   { return f->ch[0] != this && f->ch[1] != this; }
                                                                                  void change(Splay *x, int val) {
   int dir()
   { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
                                                                                     splay(x), x->val = val, x->pull();
                                                                                  int query (Splay *x, Splay *y) {
     \operatorname{ch}[d] = c;
                                                                                     split(x, y);
      if (c != \&nil) c > f = this;
                                                                                     return y->sum;
      pull();
   void give_tag(int r) {
                                                                                  3.9 KDTree [375ca2]
     if (r) swap(ch[0], ch[1]), rev = 1;
                                                                                  \begin{array}{ll} \mathbf{namespace} & \mathbf{kdt} \ \{\\ \mathbf{int} \ \mathbf{root} \ , \ \mathbf{lc} \ [\mathbf{maxn}] \ , \ \mathbf{rc} \ [\mathbf{maxn}] \ , \ \mathbf{xl} \ [\mathbf{maxn}] \ , \ \mathbf{xr} \ [\mathbf{maxn}] \ , \end{array}
   void push()
      if (ch[0] != &nil) ch[0]->give\_tag(rev);
if (ch[1] != &nil) ch[1]->give\_tag(rev);
                                                                                     yl[maxn], yr[maxn];
                                                                                  point p[maxn];
                                                                                  int build(int 1, int r, int dep = 0) {
      rev = 0;
                                                                                     if (l = r) return -1;
function<br/>bool(const point &, const point &)> f =
   void pull() {
                                                                                        [dep](const point &a, const point &b) {
  if (dep & 1) return a.x < b.x;</pre>
      // take care of the nil!
     size = ch[0]->size + ch[1]->size + 1;

sum = ch[0]->sum ^ ch[1]->sum ^ val;

if (ch[0]!= &nil) ch[0]->f = this;

if (ch[1]!= &nil) ch[1]->f = this;
                                                                                           else return a.y < b.y;
                                                                                     int m = (l + r) >> 1;
                                                                                     nth\_element(p + l, p + m, p + r, f);
} Splay::nil;
                                                                                     xl[m] = xr[m] = p[m].x;
                                                                                     yl [m] = yr [m] = p[m].y;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
                                                                                     lc[m] = build(l, m, dep + 1);
   Splay *p = x \rightarrow f;
                                                                                     if (~lc[m]) {
   int d = x->dir();
                                                                                        xl[m] = min(xl[m], xl[lc[m]]);
                                                                                         \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \end{array} 
   if (!p->isr()) p->f->setCh(x, p->dir());
   else x->f = p->f;
  p-\operatorname{setCh}(x-\operatorname{ch}[!d], d);
                                                                                        yr[m] = max(yr[m], yr[lc[m]]);
  x->setCh(p, !d);
  p->pull(), x->pull();
                                                                                     rc[m] = build(m + 1, r, dep + 1);
                                                                                          (~rc[m]) {
                                                                                        xl[m] = min(xl[m], xl[rc[m]]);
void splay (Splay *x) {
   vector<Splay*> splayVec;
                                                                                        \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{rc}[m]]);
```

yl[m] = min(yl[m], yl[rc[m]]);

```
bool erase(node *&o, int k) {
     yr[m] = max(yr[m], yr[rc[m]]);
                                                                            if (!o) return 0:
                                                                            if (o->data == k) {
    node *t = o;
  return m;
                                                                              o{-}{>}down(\,)\;,\;\;o\;=\;merge\,(\,o{-}{>}l\;,\;\;o{-}{>}r\,)\;;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
                                                                              delete t;
  if \ (q.\, x < \, x l \, [\, o\, ] \ \dot{-} \ ds \ | \ | \ \dot{q}.\, x > \, x r \, [\, o\, ] \, + \, ds \ | \ |
                                                                              return 1;
    q.y < yl[o] - ds | | q.y > yr[o] + ds
     return false;
                                                                            node *\&t = k < o->data ? o->l : o->r;
                                                                            return erase(t, k) ? o->up(), 1 : 0;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
  (a.y - b.y) * 111 * (a.y - b.y);
                                                                         void insert(node *&o, int k) {
                                                                            node *a, *b;
                                                                            split (o, a, b, k),
                                                                              o = merge(a, merge(new node(k), b));
void dfs (
                                                                         void interval(node *&o, int 1, int r) {
  const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
                                                                            node *a, *b, *c;
  long long cd = dist(p[o], q);
                                                                            split2(o, a, b, l - 1), split2(b, b, c, r);
  if (cd != 0) d = min(d, cd);
                                                                            // operate
  if ((dep & 1) & q.x < p[o].x ||
                                                                            o = merge(a, merge(b, c));
     !(dep \& 1) \& q.y < p[o].y) \{
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                         4 Flow/Matching
  } else {
                                                                         4.1 Dinic [98fb3a]
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                         {\tt struct} \  \, {\tt MaxFlow} \  \, \{ \  \, // \  \, {\tt 0-base}
                                                                            struct edge {
                                                                              int to, cap, flow, rev;
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
                                                                            vector < edge > G[MAXN];
  root = build(0, v.size());
                                                                            int s, t, dis [MAXN], cur [MAXN], n;
                                                                            int dfs(int u, int cap) {
  if (u = t || !cap) return cap;
long long nearest (const point &q) {
  long long res = 1e18;
                                                                                for \ (int \ \&i = cur[u]; \ i < (int)G[u]. \, size(); \ +\!\!+\!\! i) \ \{
  dfs(q, res, root);
                                                                                 edge &e = G[u][i]
  return res;
                                                                                 if (dis[e.to] = dis[u] + 1 \&\& e.flow != e.cap) {
                                                                                    int df = dfs(e.to, min(e.cap - e.flow, cap));
} // namespace kdt
                                                                                    if (df) {
                                                                                      e.flow += df;
G[e.to][e.rev].flow -= df;
3.10 Treap [5ab1a1]
                                                                                      return df;
struct node {
  int data, sz;
node *l, *r;
                                                                                   }
                                                                                }
  node({\tt int}\ k)\ :\ data(k)\,,\ sz(1)\,,\ l(0)\,,\ r(0)\ \{\}
  void up() {
                                                                               dis[u] = -1;
                                                                              {\tt return} \ 0;
     sz = 1;
     if (1) sz += 1->sz;
    if (r) sz += r->sz;
                                                                            bool bfs() {
                                                                              fill_n(dis, n, -1);
  void down() {}
                                                                              queue<int> q;
                                                                              q.push(s), dis[s] = 0;
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
   if (rand() % (sz(a) + sz(b)) < sz(a))</pre>
                                                                               while (!q.empty()) {
                                                                                 int tmp = q.front();
                                                                                 q.pop();
                                                                                 for (auto &u : G[tmp])
     if (!~dis[u.to] && u.flow != u.cap) {
                                                                                      q.push(u.to);
dis[u.to] = dis[tmp] + 1;
  return b->down(), b->l = merge(a, b->l), b->up(), b;
void split (node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
                                                                              return dis[t] != -1;
  o->down();
  if (o->data \le k)
                                                                            int maxflow(int _s, int _t) {
    a = o, split(o->r, a->r, b, k), a->up();
                                                                              s = _s, t = _t;
  else b = o, split(o->l, a, b->l, k), <math>b->up();
                                                                               int flow = 0, df;
                                                                              while (bfs()) {
                                                                                 \begin{array}{l} \text{fill\_n}(\text{cur}, n, 0); \\ \text{while} ((\text{df} = \text{dfs}(\text{s}, \text{INF}))) \text{ flow} += \text{df}; \end{array}
void split2 (node *o, node *&a, node *&b, int k) {
 if (sz(o) \le k) return a = o, b = 0, void();
  o > down();
  if (sz(o->1) + 1 \le k)
                                                                              return flow;
  a = o, split2(o->r, a->r, b, k - sz(o->l) - 1);
else <math>b = o, split2(o->l, a, b->l, k);
                                                                            void init(int _n) {
  o > up();
                                                                              n = \underline{n};
                                                                              for (int i = 0; i < n; ++i) G[i].clear();
node *kth(node *o, int k) {
  if (k \le sz(o->l)) return kth(o->l, k);
                                                                            void reset() {
  if (k = sz(o->l) + 1) return o;
                                                                              for (int i = 0; i < n; ++i)
  return kth(o->r, k-sz(o->l)-1);
                                                                                 for (auto \& j : G[i]) j.flow = 0;
int Rank(node *o, int key) {
                                                                            void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, (int)G[v].size()});
  G[v].pb(edge{u, 0, 0, (int)G[u].size() - 1});
  if (!o) return 0;
  if (o->data < key)
     return sz(o->l) + 1 + Rank(o->r, key);
  else return Rank(o->l, key);
                                                                        };
```

4.2 Bipartite Matching* [784535]

```
struct Bipartite_Matching { // 0-base
  vector < int > G[N + 1];
  bool dfs(int u) {
    int e = G[u][i];
      if (mq[e] == 1
           || (dis[mq[e]] = dis[u] + 1 \& dfs(mq[e]))|
        return mp[mq[e] = u] = e, 1;
    return dis[u] = -1, 0;
  bool bfs() {
    queue<int> q;
    fill_n (dis, l + 1, -1);
for (int i = 0; i < l; ++i)
      if (!~mp[i])
        q.push(i), dis[i] = 0;
    while (!q.empty())
      int u = q.front();
      q.pop();
      for (int e : G[u])
        if (!~dis[mq[e]])
          q.push(mq[e]), dis[mq[e]] = dis[u] + 1;
    return dis[1] != -1;
  int matching() {
    int res = 0;
    fill\_n\left(mp,\ l\ ,\ -1\right),\ fill\_n\left(mq,\ r\ ,\ l\ \right);
    while (bfs()) {
      fill_n(cur, 1, 0);

for(int i = 0; i < 1; ++i)
        res += (! \sim mp[i] \&\& dfs(i));
    return res; // (i, mp[i] != -1)
  void add_edge(int s, int t) { G[s].pb(t); }
  void init(int _l, int _r) {
    l = _l, r = _r;
    for (int i = 0; i \le l; ++i)
      G[i].clear();
};
```

4.3 Kuhn Munkres* [4b3863]

```
struct KM { // 0-base, maximum matching
 11 w[N] [N], h1 [N], hr [N], slk [N]; int f1 [N], fr [N], pre [N], qu [N], ql, qr, n; bool v1 [N], vr [N];
  void init(int _n) {
   void add_edge(int a, int b, ll wei) {
   w[a][b] = wei;
  bool Check(int x) {
if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0:
  void bfs(int s) {
    fill_n (slk
        , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) 
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
         else slk[x] -= d;
        if (vr[x]) hr[x] = d;
```

4.4 MincostMaxflow* [1c78db]

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
     ll from, to, cap, flow, cost, rev;
   } *past[N];
   vector < Edge > G[N];
  \begin{array}{lll} & \text{int inq}\left[N\right], \ n, \ s, \ t; \\ & \text{ll dis}\left[N\right], \ up\left[N\right], \ pot\left[N\right]; \\ & \text{bool BellmanFord}\left(\right) \ \left\{ \end{array} \right.
     fill_n(dis, n, INF), fill_n(inq, n, 0);
     queue<int> q;
     if (!inq[u]) inq[u] = 1, q.push(u);
       }
     };
     relax(s, 0, INF, 0);
while (!q.empty()) {
       int u = q.front();
        q.pop(), inq[u] = 0;
        for (auto &e : G[u]) {
          11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
                (e.to, d2, min(up[u], e.cap - e.flow), &e);
       }
     return dis[t] != INF;
  void solve (int
       , int _t, ll &flow, ll &cost, bool neg = true) {
     s = _s, t = _t, flow = 0, cost = 0;
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     for (; BellmanFord(); copy_n(dis, n, pot)) {
        for (int
        i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
        for (int i = t; past[i]; i = past[i]->from) {
          auto &e = *past[i]
           e.flow += up[t], G[e.to][e.rev].flow -= up[t];
     }
  }
  void init(int _n) {
     n = \underline{n}, fill_n(pot, n, 0);
     for (int i = 0; i < n; ++i) G[i].clear();
  void add_edge(ll a, ll b, ll cap, ll cost) {
   G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
   G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
```

4.5 Maximum Simple Graph Matching* [0fe1c3]

```
struct Matching { // 0-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
        if (vis[x] == tk) return x;
        vis[x] = tk;
        x = Find(pre[match[x]]);
    }
  }
  void Blossom(int x, int y, int l) {
```

```
for (; Find(x) != 1; x = pre[y]) {
    pre[x] = y, y = match[x];
if (s[y] = 1) q.push(y), s[y] = 0;
    for (int z: \{x, y\}) if (fa[z] = z) fa[z] = 1;
 }
bool Bfs(int r) {
  iota(ALL(fa), 0); fill(ALL(s), -1);
  q = queue<int>(); q.push(r); s[r] = 0;
for (; !q.empty(); q.pop()) {
    for (int x = q. front(); int u : G[x])
      if(s[u] = -1) {
         last =
                 match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]); s[match[u]] = 0;
      } else if (!s[u] \&\& Find(u) != Find(x)) {

int l = LCA(u, x);

         Blossom(x, u, l); Blossom(u, x, l);
  return false;
Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis
(n+1), \text{ pre}(n+1, n), \text{ match}(n+1, n), G(n) \{\}
\text{void add\_edge}(\text{int } u, \text{ int } v)
{ G[u].pb(v), G[v].pb(u); }
int solve() {
  int ans = 0;
  for (int x = 0; x < n; ++x)
     if (match[x] == n) ans += Bfs(x); 
  return ans:
} // match[x] == n means not matched
```

4.6 Maximum Weight Matching* [1ec446]

```
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
      struct edge { int u, v, w; }; int n, nx;
       vector<int> lab; vector<vector<edge>>> g;
       \begin{array}{l} \text{vector} < \text{int} > \text{slk} \;,\;\; \text{match} \;,\; \text{st} \;,\; \text{pa} \;,\; S \;,\; \text{vis} \;; \\ \text{vector} < \text{vector} < \text{int} > \; \text{flo} \;,\;\; \text{flo\_from} \;;\;\; \text{queue} < \text{int} > \; q \;; \\ \text{WeightGraph} (\;\; \text{int} \;\; \text{n}\_) \;:\;\; \text{n}(\text{n}\_) \;,\;\; \text{nx}(\text{n} \;\;^* \;\; 2) \;,\;\; \text{lab}(\text{nx} \;+\; 1) \;,\;\; \\ \text{g}(\text{nx} \;+\; 1 \;,\;\; \text{vector} < \text{edge} > (\text{nx} \;+\; 1) \;,\;\; \text{slk} (\text{nx} \;+\; 1) \;,\;\; \\ \end{array} 
             flo (nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
match = st = pa = S = vis = slk;
            REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
       int E(edge e)
      { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; } void update_slk(int u, int x, int &s) { if (!s || E(g[u][x]) < E(g[s][x])) s = u; }
       void set_slk(int x) {
               slk[x] = 0;
                      \begin{tabular}{ll} \be
                            update\_slk(u, x, slk[x]);
       void q_push(int x)
              if (x \le n) q.push(x);
              else for (int y : flo[x]) q_push(y);
       void set_st(int x, int b) {
              st[x] = b;
              if (x > n) for (int y : flo[x]) set_st(y, b);
       vector<int> split_flo(auto &f, int xr) {
              auto it = find(ALL(f), xr);
              if (auto pr = it - f.begin(); pr % 2 == 1)
                     reverse(1 + ALL(f)), it = f.end() - pr;
              auto res = vector(f.begin(), it);
              return f.erase(f.begin(), it), res;
       void set_match(int u, int v) {
            match[u] = g[u][v].v;
              if (u \le n) return;
              \begin{array}{ll} \textbf{int} & \textbf{xr} = \textbf{flo\_from} \, [\, \textbf{u} \,] \, [\, \textbf{g} \, [\, \textbf{u} \,] \, [\, \textbf{v} \,] \, . \, \textbf{u} \,] \,; \end{array}
            set_match(xr, v); f.insert(f.end(), ALL(z));
       void augment(int u, int v) {
```

```
for (;;) {
    int xnv = st[match[u]]; set\_match(u, v);
     if (!xnv) return;
    set_{match}(v = xnv, u = st[pa[xnv]]);
int lca(int u, int v) {
  static int t = 0; ++t;
  for (++t; u \mid | v; swap(u, v)) if (u) {
    if (vis[u] == t) return u;
     vis[u] = t, u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0, match[b] = match[o];
  vector < int > f = \{o\};
  for (int t : {u, v}) {
     reverse(1 + ALL(f));
     f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set\_st(b, b);
  REP(x, 1, nx) g[b][x].w = g[x][b].w = 0; fill (ALL(flo_from[b]), 0);
  for (int xs : flo[b]) {
    REP(x, 1, nx)
       if (g[b][x].w = 0 \mid | E(g[xs][x]) < E(g[b][x]))
         g[b][x] = g[xs][x], g[x][b] = g[x][xs];
             1, n
       if (flo_from [xs][x]) flo_from [b][x] = xs;
  set_slk(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x: split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
     slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
  for (int x : flo[b])

if (x = xr) S[x] = 1, pa[x] = pa[b];
     else S[x] = -1, set\_slk(x);
  \operatorname{st}[b] = 0;
bool on_found_edge(const edge &e) {
   \begin{array}{lll} & \text{if (int } u = st \, [e \cdot u] \;, \; v = st \, [e \cdot v] \;; \; S[v] = -1) \; \{ \\ & \text{int } nu = st \, [match[v]] \;; \; pa[v] = e \cdot u \;; \; S[v] = 1 \;; \\ & \text{slk} \, [v] = slk \, [nu] = S[nu] = 0 \;; \; q\_push(nu) \;; \\ \end{array} 
  else if (S[v] = 0)
     if (int o = lca(u, v)) add_blossom(u, o, v);
     else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
  fill(ALL(S), -1), fill(ALL(slk), 0);
    = queue < int > ();
  for (;;)
     while (SZ(q)) {
       int u = q.front(); q.pop();
        if (S[st[u]] = 1) continue;
       REP(\,v\,,\ 1\,,\ n\,)
          update\_slk(u, st[v], slk[st[v]]);
            else if
                   (on_found_edge(g[u][v])) return true;
     int d = INF;
     \begin{array}{lll} REP(b, \ n+\stackrel{'}{1}, \ nx) & if \ (st[b] = b \&\& \ S[b] == 1) \\ d = min(d, \ lab[b] \ / \ 2); \end{array} 
    REP(x, 1, nx)
       if (int
             s = slk[x]; st[x] == x \&\& s \&\& S[x] <= 0
         d = min(d, E(g[s][x]) / (S[x] + 2));
    REP(u, 1, n)
       if (S[st[u]] == 1) lab[u] += d;
```

void init(int _n) {

 $n = _n;$ for (int i = 0; i < n + 2; ++i)

G[i].clear(), cnt[i] = 0;

```
\begin{array}{ll} \text{else if } (S[st[u]] =\!\!\!\!= 0) \ \{ \\ \text{if } (lab[u] <\!\!\!= d) \ \text{return false}; \end{array}
                                                                                    \begin{array}{l} void \ add\_edge(int \ u, \ int \ v, \ int \ lcap \,, \ int \ rcap) \ \{ \\ cnt[u] \ -= \ lcap \,, \ cnt[v] \ += \ lcap \,; \end{array} 
                                                                                      G[\,u\,] \cdot pb\big(\,edge\big\{v\,,\ rcap\,,\ lcap\,,\ SZ(G[\,v\,]\,)\,\big\}\big)\,;
             lab[u] -= d;
                                                                                      G[v].pb(edge\{u, 0, 0, SZ(G[u]) -
        void add_edge(int u, int v, int cap) {
   G[u].pb(edge{v, cap, 0, SZ(G[v])});
        REP(x, 1, nx)
           if (int s = slk[x]; st[x] == x &&
                                                                                      G[v].pb(edge\{u, 0, 0, SZ(G[u]) - 1\});
                s \&\& st[s] != x \&\& E(g[s][x]) == 0
              if (on_found_edge(g[s][x])) return true;
                                                                                   int dfs(int u, int cap) {
                                                                                       \begin{array}{lll} & \text{if } (u = t \mid \mid ! \operatorname{cap}) \text{ return } \operatorname{cap}; \\ & \text{for } (\operatorname{int \&i} = \operatorname{cur}[u]; \text{ } i < \operatorname{SZ}(G[u]); \text{ } +\!\!+\!\! i) \end{array} \} 
        REP(b, n + 1, nx)

if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                                         edge &e = G[u][i];
if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
              expand_blossom(b);
                                                                                            int df = dfs(e.to, min(e.cap - e.flow, cap));
     return false;
                                                                                            if (df) {
   }
   pair<ll, int> solve() {
                                                                                               e.flow += df, G[e.to][e.rev].flow -= df;
     fill(ALL(match), 0);
                                                                                               return df;
     REP(u, 0, n) st [u] = u, flo [u]. clear();
                                                                                        }
     int w_max = 0;
     dis[u] = -1;
        w_{max} = max(w_{max}, g[u][v].w);
                                                                                      return 0;
     fill (ALL(lab), w_max);
int n_matches = 0; 11 tot_weight = 0;
                                                                                   bool bfs() {
                                                                                      fill_n(dis, n + 3, -1);
      while (matching()) ++n_matches;
                                                                                      queue<int> q;
     REP(u, 1, n) \ if \ (match[u] \ \&\& \ match[u] < u)
                                                                                      q.push(s), dis[s] = 0;
        tot\_weight += g[u][match[u]].w;
                                                                                      while (!q.empty()) {
                                                                                         int \dot{\mathbf{u}} = \mathbf{q}. \, \text{front}();
     return make_pair(tot_weight, n_matches);
                                                                                         q.pop();
                                                                                         for (edge &e : G[u])
   void add_edge(int u, int v, int w)
   \{g[u][v].w = g[v][u].w = w; \}
                                                                                            if (!~dis[e.to] & e.flow != e.cap)
                                                                                              q.push(e.to), dis[e.to] = dis[u] + 1;
        SW-mincut [6621f5]
4.7
                                                                                      return dis[t] != -1;
int maxflow(int _s, int _t) {
                                                                                      s = \underline{s}, t = \underline{t};

int flow = 0, df;
   \label{eq:main_section}  \begin{array}{ll} \text{int} & \text{vst} \left[ \text{MXN} \right] \;, \;\; \text{edge} \left[ \text{MXN} \right] \left[ \text{MXN} \right] \;, \;\; \text{wei} \left[ \text{MXN} \right] ; \end{array}
                                                                                      while (bfs()) {
   void init(int n)
                                                                                         fill_n(cur, n + 3, 0);
     REP fill_n (edge[i], n, 0);
                                                                                         while ((df = dfs(s, INF))) flow += df;
  void addEdge(int u, int v, int w){
  edge[u][v] += w; edge[v][u] += w;
                                                                                      return flow;
                                                                                   bool solve() {
   int search(int &s, int &t, int n){
                                                                                      int sum = 0;
     fill_n(vst, n, 0), fill_n(wei, n, 0);
                                                                                      for (int i = 0; i < n; ++i)
     s = t = -1;
                                                                                         if (cnt[i] > 0)
     int mx, cur;
                                                                                         \begin{array}{lll} & \text{add\_edge}(n+1,\ i\ ,\ cnt[i])\ ,\ sum\ +=\ cnt[i];\\ & \text{else} & \text{if}\ (cnt[i]<0)\ add\_edge(i\ ,\ n+2,\ -cnt[i])\ ; \end{array}
     for (int j = 0; j < n; +++j) {
        mx = -1, cur = 0;
                                                                                      if (\text{sum }!=\text{maxflow}(n+1, n+2)) \text{ sum }=-1;
        REP if (wei[i] > mx) cur = i, mx = wei[i];
                                                                                      for (int i = 0; i < n; ++i)
        vst[cur] = 1, wei[cur] = -1;
                                                                                         if (cnt[i] > 0)
        s = t; t = cur
                                                                                           G[n + 1].pop_back(), G[i].pop_back();
lse if (cnt[i] < 0)
        REP if (!vst[i]) wei[i] += edge[cur][i];
                                                                                           G[i].pop\_back(), G[n + 2].pop\_back();
     return mx;
                                                                                      return sum != -1;
   int solve(int n) {
                                                                                   int solve(int _s, int
     int res = INF;
                                                                                      add\_edge(\_t, \_s, INF);
     for (int x, y; n > 1; n--){
                                                                                      if (!solve()) return -1; // invalid flow
        res = min(res, search(x, y, n));
REP edge[i][x] = (edge[x][i] += edge[y][i]);
                                                                                      int x = G[\_t].back().flow;
                                                                                      return G[_t].pop_back(), G[_s].pop_back(), x;
        REP {
           edge[y][i] = edge[n - 1][i];
edge[i][y] = edge[i][n - 1];
                                                                                };
        4.9 Gomory Hu tree* [11be99]
                                                                                MaxFlow Dinic;
     return res;
                                                                                 int g [MAXN];
} sw;
                                                                                 void GomoryHu(int n) { // 0-base
                                                                                   fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
        BoundedFlow*(Dinic*) [4a793f]
                                                                                      Dinic.reset();
struct BoundedFlow { // 0-base
                                                                                      \begin{array}{lll} add\_edge(i\;,\;g[\,i\,]\;,\;Dinic.maxflow(\,i\;,\;g[\,i\,])\,)\,;\\ for\;(\,int\;\;j=\,i\;+\;1;\;\;j<=\,n;\;+\!\!+\!\!j\,) \end{array}
   struct edge {
     int to, cap, flow, rev;
                                                                                          if (g[j] = g[i] &  \text{ $\sim$ Dinic.dis}[j] ) 
                                                                                           g\,[\,j\,] \;=\; i\;;
  vector<edge> G[N];
int n, s, t, dis[N], cur[N], cnt[N];
```

struct MinCostCirculation { // 0-base
 struct Edge {

4.10 Minimum Cost Circulation* [ba97cf]

```
ll from, to, cap, fcap, flow, cost, rev; } *past[N];
   vector < Edge > G[N];
11 dis [N], inq [N], n;
   void BellmanFord(int s) {
      fill\_n \left( \, dis \, , \, \, n \, , \, \, INF \right) \, , \  \, fill\_n \left( \, inq \, , \, \, n \, , \, \, \, 0 \right) ;
      queue<int> q;
      auto relax = [&](int u, ll d, Edge *e) {
  if (dis[u] > d) {
           dis[u] = d, past[u] = e;
            if (!inq[u]) inq[u] = 1, q.push(u);
        }
      relax(s, 0, 0);
      while (!q.empty()) {
        int u = q.front();
         q.pop(), inq[u] = 0;
         for (auto &e : G[u])
            if (e.cap > e.flow)
              relax(e.to, dis[u] + e.cost, &e);
   }
   void try_edge(Edge &cur) {
      if (cur.cap > cur.flow) return ++cur.cap, void();
      BellmanFord(cur.to);
      if (dis[cur.from] + cur.cost < 0) {
    ++cur.flow, --G[cur.to][cur.rev].flow;</pre>
         for (int
           \begin{tabular}{ll} $i=cur.from; $past[i]$; $i=past[i]$->from) { $auto &e=*past[i]$; $++e.flow, $--G[e.to][e.rev].flow; $} \end{tabular}
        }
     ++cur.cap;
   void solve(int mxlg) {
  for (int b = mxlg; b >= 0; --b) {
         for (int i = 0; i < n; ++i)
           for (auto &e : G[i])
        e cap *= 2, e flow *= 2;
for (int i = 0; i < n; ++i)
for (auto &e : G[i])
              if (e.fcap >> b & 1)
                 try_edge(e);
     }
   void init (int \underline{\phantom{a}}n) { n = \underline{\phantom{a}}n;
      for (int i = 0; i < n; ++i) G[i].clear();
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].pb(Edge
            \{a, b, 0, cap, 0, cost, SZ(G[b]) + (a = b)\};
     G[b].pb(Edge\{b, a, 0, 0, -cost, SZ(G[a]) - 1\});
} mcmf; // O(VE * ElogC)
```

4.11 Flow Models

- $\bullet \quad Maximum/Minimum flow with lower bound/Circulation problem\\$
 - 1. Construct supersource S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite $\operatorname{graph}(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct supersource S and $\sinh T$
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1

- 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) = (0, d(v))
- 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer T
- $2. \ \ {\it Constructa\, max\, flow\, model}, {\it let}\, K\, {\it be\, the \, sum\, of \, all\, weights}$
- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - $3. \ \ The {\it mincut} is equivalent to the {\it maximum} profit of a subset of projects.$
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
 - 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \ge 0$$

4.12 matching

- 最大匹配+最小邊覆蓋=V
- 最大獨立集+最小點覆蓋=V
- 最大匹配=最小點覆蓋
- 最小路徑覆蓋數=V-最大匹配數

5 String

5.1 KMP [5a0728]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
  }
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  }
  return ans;
}</pre>
```

5.2 Z-value* [b47c17]

5.3 Manacher* [lad8ef]

```
int z[MAXN]; // 0-base
/* center i: radius z[i * 2 + 1] / 2
    center i, i + 1: radius z[i * 2 + 2] / 2
    both aba, abba have radius 2 */
void Manacher(string tmp) {
    string s = "%";
    int l = 0, r = 0;
    for (char c : tmp) s.pb(c), s.pb("%");
    for (int i = 0; i < SZ(s); ++i) {
        z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
        while (i - z[i] >= 0 && i + z[i] < SZ(s)
        && s[i + z[i]] == s[i - z[i]]) ++z[i];
}</pre>
```

int R = q.front();

q.pop(), ord[t++] = R;

```
\quad \text{for (int } i = 0; \ i < sigma; ++i)
      if (z[i] + i > r) r = z[i] + i, l = i;
                                                                                        if (~nx[R][i])
}
                                                                                           \begin{array}{lll} & \text{int } X = rnx \left[ R \right] \left[ \ i \ \right] = nx \left[ R \right] \left[ \ i \ \right], \ Z = fl \left[ R \right]; \\ & \text{for } (; \ Z \&\& \ ! \sim nx \left[ Z \right] \left[ \ i \ \right]; \ ) \ Z = fl \left[ Z \right]; \\ & fl \left[ X \right] = Z \ ? \ nx \left[ Z \right] \left[ \ i \ \right] \ : \ 1, \ q.push (X); \end{array}
 5.4 SAIS* [6f26bc]
auto sais(const auto &s) {
                                                                                        else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
   const int n = SZ(s), z = ranges :: max(s) + 1;
                                                                                  }
   if (n = 1) return vector{0};
   vector < int > c(z); for (int x : s) + c[x];
                                                                                void solve() {
  for (int i = top - 2; i > 0; --i)
    cnt[fl[ord[i]]] += cnt[ord[i]];
   partial\_sum(ALL(c), begin(c));
   } ac;
      t[i] = (
          s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
                                                                             5.6 Smallest Rotation [4f469f]
   auto is_lms = views::filter([&t](int x) {
      return x && t[x] && !t[x - 1];
                                                                             string mcp(string s) {
                                                                                int n = SZ(s), i = 0, j = 1;
   auto induce = [&] {
                                                                                s += s;
     for (auto x = c; int y : sa)
                                                                                while (i < n \& j < n) {
      if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
for (auto x = c; int y : sa | views::reverse)
if (y--) if (t[y]) sa[--x[s[y]]] = y;
                                                                                   int k = 0;
                                                                                   while (k < n \& s[i + k] = s[j + k]) ++k;
                                                                                   if (s[i + k] \le s[j + k]) j += k + 1;
                                                                                   else i += k + 1;
   vector<int> lms, q(n); lms.reserve(n);
                                                                                   if (i == j) ++j;
   for (auto x = c; int i : I | is_lms)

q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
                                                                                int ans = i < n ? i : j;
   induce(); vector < int > ns(SZ(lms));
                                                                                return s.substr(ans, n);
   for (int j = -1, nz = 0; int i : sa \mid is\_lms) {
     if (j >= 0) {
        5.7 De Bruijn sequence* [a09470]
         ns[q[i]] = nz += lexicographical\_compare(
                                                                             constexpr int MAXC = 10, MAXN = 1e5 + 10;
             \begin{array}{l} \text{begin}(s) + j, & \text{begin}(s) + j + \text{len}, \\ \text{begin}(s) + i, & \text{begin}(s) + i + \text{len}); \end{array}
                                                                             struct DBSeq {
                                                                                int C, N, K, L, buf [MAXC * MAXN]; // K <= \mathbb{C}^N
                                                                                void dfs(int *out, int t, int p, int &ptr) {
      j \ = \ i \ ;
                                                                                   if (ptr >= L) return;
                                                                                   if (t > N) {
if (N \% p) return;
   fill(ALL(sa), 0); auto nsa = sais(ns);
   for (auto x = c; int y : nsa \mid views::reverse)

y = lms[y], sa[--x[s[y]]] = y;
                                                                                      for (int i = 1; i \le p \&\& ptr < L; ++i)
                                                                                        \operatorname{out}[\operatorname{ptr}++] = \operatorname{buf}[i];
   return induce(), sa;
                                                                                   } else
                                                                                     \begin{array}{l} buf[t] = buf[t - p], \ dfs(out, \ t + 1, \ p, \ ptr); \\ for \ (int \ j = buf[t - p] + 1; \ j < C; ++j) \end{array}
 // sa[i]: sa[i]-th suffix
 is the i-th lexicographically smallest suffix. // hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
                                                                                        buf[t] = j, dfs(out, t + 1, t, ptr);
 struct Suffix {
   int n; vector<int> sa, hi, ra;
                                                                                void solve(int _c, int _n, int _k, int *out) {
   Suffix
     \begin{array}{l} (const\ auto\ \&\_s,\ int\ \_n):\ n(\_n)\,,\ hi\,(n)\,,\ ra\,(n)\ \{\\ vector<int>\ s(n+1);\ //\ s\,[n]\ =\ 0;\\ copy\_n(\_s,\ n,\ begin\,(s));\ //\ \_s\ shouldn\ 't\ contain\ 0 \end{array}
                                                                                  int p = 0;
                                                                                  sa = sais(s); sa.erase(sa.begin())
      for (int i = 0; i < n; ++i) ra[sa[i]] = i;
      for (int i = 0, h = 0; i < n; ++i) {
  if (!ra[i]) { h = 0; continue; }
                                                                             } dbs;
                                                                             5.8 Extended SAM* [64c3b7]
         for (int j = sa[ra[i] - 1]; max
                   j) + h < n \& s[i + h] = s[j + h];) + h;
                                                                             hi[ra[i]] = h ? h-- : 0;
   }
};
                                                                                int cnt[N * 2]; // occurence
        Aho-Corasick Automatan* [794a77]
                                                                                int newnode()
                                                                                   fill_n (next[tot], CNUM, 0);
 struct AC_Automatan {
                                                                                   len[tot] = cnt[tot] = link[tot] = 0;
   int nx[len][sigma], fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
                                                                                   return tot++;
                                                                                void init() { tot = 0, newnode(), link[0] = -1; }
   int newnode() {
      fill_n(nx[top], sigma, -1);
                                                                                int insertSAM(int last, int c) {
                                                                                  int cur = next[last][c];
len[cur] = len[last] + 1;
      return top++;
                                                                                   int p = link[last];
   void init() \{ top = 1, newnode(); \}
   int input(string &s) {
                                                                                   while (p != -1 && ! next[p][c])
                                                                                   next[p][c] = cur, p = link[p];
if (p = -1) return link[cur] = 0, cur;
      int X = 1;
      for (char c : s) {    if (!\simnx[X][c - 'A']) nx[X][c - 'A'] = newnode();    X = nx[X][c - 'A'];
                                                                                   int q = next[p][c];
                                                                                   if (len
                                                                                        [p] + 1 = len[q] return link[cur] = q, cur;
      return X; // return the end node of string
                                                                                   int clone = newnode();
                                                                                   for (int i = 0; i < ONUM; ++i)
   void make_fl() {
                                                                                     next[
      queue<int> q;
                                                                                           clone][i] = len[next[q][i]] ? next[q][i] : 0;
      q.push(1), fl[1] = 0;
                                                                                   len[clone] = len[p] + 1;
      for (int t = 0; !q.empty(); ) {
                                                                                   while (p != -1 && next[p][c] == q)
```

next[p][c] = clone, p = link[p];

link[link[cur] = clone] = link[q];

```
link[q] = clone;
                                                                     void main_lorentz(const string &s, int sft = 0) {
     return cur;
                                                                       const int n = s.size();
                                                                       if (n = 1) return;
const int nu = n / 2, nv = n - nu;
   void insert (const string &s) {
     int cur = 0;
                                                                       const string u = s.substr(0, nu), v = s.substr(nu),
     for (auto ch : s) {
                                                                              ru(u.rbegin
       int &nxt = next[cur][int(ch - 'a')];
                                                                                   (), u.rend()), rv(v.rbegin(), v.rend());
                                                                       \begin{array}{ll} main\_lorentz(u, sft), \; main\_lorentz(v, sft + nu); \\ const \; auto \; z1 = Zalgo(ru), \; z2 = Zalgo(v + '\#' + u) \end{array}
        if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
                                                                                   z3 = Zalgo(ru' + '\#' + rv), z4 = Zalgo(v);
     }
                                                                       auto get_z = [](const vector<int>&z, int i) {
   }
   void build() {
                                                                         return
     queue < int > q;
                                                                               (0 \le i \text{ and } i \le (int)z.size()) ? z[i] : 0; };
                                                                       auto add_rep
     q.push(0);
     while (!q.empty()) {
                                                                             = [&](bool left, int c, int l, int k1, int k2) {
       int cur = q.front();
                                                                         const
                                                                               int L = max(1, l - k2), R = min(l - left, k1);
        q.pop();
        for (int i = 0; i < CNUM; ++i)
  if (next[cur][i])</pre>
                                                                          if (L > R) return;
                                                                         if (left)
            q.push(insertSAM(cur, i));
                                                                               rep[l].emplace\_back(sft + c - R, sft + c - L);
                                                                          else rep[1].emplace_back
                                                                              (sft + c - R - l + 1, sft + c - L - l + 1);
     vector < int > lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];
     partial_sum(ALL(lc), lc.begin());
                                                                       for (int cntr = 0; cntr < n; cntr++) {
                                                                         int 1, k1, k2;
     for (int i
          = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
                                                                         if (cntr < nu) {
                                                                            l = nu - cntr;
   void solve() {
                                                                            k1 = get_z(z1, nu - cntr);
     for (int i = tot - 2; i >= 0; --i)
                                                                            k2 = get_z(z2, nv + 1 + cntr);
        cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
                                                                         } else {
                                                                            l = cntr - nu + 1;
 };
                                                                            k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
                                                                            k2 = get_z(z4, (cntr - nu) + 1);
5.9 PalTree* [d7d2cf]
                                                                          if (k1 + k2 >= 1)
struct palindromic_tree {
                                                                            add\_rep\,(\,cntr\,<\,nu\,,\ cntr\,,\ l\,,\ k1\,,\ k2\,)\,;
   struct node {
     int next[26], fail, len;
                                                                    int cnt, num; // cnt: appear times, num: number of
                      // pal. suf.
                                                                     6
                                                                          Math
     node(int \ l = 0)': fail(0), len(1), cnt(0), num(0) {
                                                                     6.1 ax+by=gcd(only exgcd *) [7b833d]
       for (int i = 0; i < 26; ++i) next[i] = 0;
                                                                     pll exgcd(ll a, ll b)
                                                                       if (b = 0) return pll(1, 0);
   vector<node> St;
                                                                       ll p = a / b;
   vector<char> s;
                                                                       pll q = exgcd(b, a \% b);
   int last, n;
                                                                       palindromic\_tree() : St(2), last(1), n(0) 
     St[0].fail = 1, St[1].len = -1, s.pb(-1);
                                                                     /* ax+by=res, let x be minimum non-negative
                                                                    g, p = gcd(a, b), exgcd(a, b) * res / g
   inline void clear() {
                                                                     \begin{array}{l} \text{if } p.X < 0: \ t = (abs(p.X) + b \ / \ g - 1) \ / \ (b \ / \ g) \\ \text{else: } t = -(p.X \ / \ (b \ / \ g)) \\ p += (b \ / \ g, \ -a \ / \ g) \ * \ t \ */ \\ \end{array} 
     St.clear(), s.clear(), last = 1, n = 0; St.pb(0), St.pb(-1);
     St[0]. fail = 1, s.pb(-1);
                                                                     6.2 Floor and Ceil [692c04]
   int floor (int a, int b)
                                                                    { return \hat{a} / \hat{b} - (a % \hat{b} && (a < 0) \hat{} (b < 0)); } int ceil(int a, int b)
       x = St[x]. fail;
     return x;
                                                                    { return a / b + (a \% b \&\& (a < 0) \cap (b > 0)); }
   inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
  int cur = get_fail(last);
                                                                     6.3 Floor Enumeration [7cbcdf]
                                                                      / enumerating x = floor(n / i), [l,
     if (!St[cur].next[c]) {
  int now = SZ(St);
                                                                    for (int l = 1, r; l <= n; l = r + 1) {
int x = n / l;
        St.pb(St[cur].len + 2);
                                                                       r = n / x;
        St[now].fail =
                                                                    }
          St[get_fail(St[cur].fail)].next[c];
        St[cur].next[c] = now;
                                                                     6.4 Mod Min [9118e1]
        St[now].num = St[St[now].fail].num + 1;
                                                                    // \min\{k \mid 1 \le ((ak) \mod m) \le r\}, no solution -> -1 ll \max_{\min}(ll \ a, \ ll \ m, \ ll \ l, \ ll \ r)  { if (a == 0) return l \ ? \ -1 : \ 0; if (ll \ k = (l + a \ -1) \ / \ a; \ k * a <= r)
     last = St[cur].next[c], ++St[last].cnt;
   inline void count() { // counting cnt
     auto i = St.rbegin();
                                                                         return k;
     for (; i != St.rend(); ++i) {
                                                                       ll\ b = m\ /\ a\,,\ c = m\ \%\ a\,;
       St[i->fail].cnt += i->cnt;
                                                                       return (1 + y * c + a - 1) / a + y * b;
   inline int size() { // The number of diff. pal.
return SZ(St) - 2;
                                                                     6.5 Linear Mod Inverse [5a4cbf]
 };
                                                                    inv[1] = 1:
                                                                     for ( int i = 2; i
 5.10 Main Lorentz [615b8f]
                                                                          \langle = N; ++i \rangle inv[i] = ((mod-mod/i)*inv[mod%i])%mod;
| \text{vector} < \text{pair} < \text{int}, \text{ int} \gg \text{rep}[kN]; // 0-\text{base}[l, r]
```

6.6 Linear Filter Mu [ac2ac3]

6.7 Gaussian integer gcd [763e59]

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
     (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
     ll c = a.real() * b.real() + a.imag() * b.imag();
     ll d = a.imag() * b.real() - a.real() * b.imag();
     ll r = b.real() * b.real() + b.imag() * b.imag();
     if (c % r == 0 && d % r == 0) return b;
     return gaussian_gcd
          (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
}
```

6.8 GaussElimination [6308be]

```
\begin{tabular}{ll} \begin{tabular}{ll} void $GAS(V<V<double>>&vc) & \{ \end{tabular}
      int len = vc.size();
      for (int i = 0; i < len; ++i)
             int idx = find_if(vc.begin()+i, vc.end(), [&](
                   if ( idx = len ) continue
            if( i != idx ) swap( vc[idx], vc[i] );
double pivot = vc[i][i];
            for_each( vc[i].begin(), vc
            iol_each( vc[i].bcgin(), vc
  [i].end(), [&]( auto &a ) { a/=pivot; } );
for( int j = 0; j < len; ++j ) {
  if( i == j ) continue;
  if( vc[j][i]!= 0 ) {
    double mul = vc[j][i]/vc[i][i];
    trypeform( va[i] bosin() va[i] end</pre>
                         transform(vc[j].begin(), vc[j].end
                                (), vc[i].begin(), vc[j].begin(),
                                     [&](auto &a, auto &b) {
                                      return a-b*mul;
                                     });
                   }
            }
      }
};
```

6.9 floor sum* [49de67]

```
11 floor_sum(ll n, ll m, ll a, ll b) {
    ll ans = 0;
    if (a >= m)
        ans += (n - 1) * n * (a / m) / 2, a %= m;
    if (b >= m)
        ans += n * (b / m), b %= m;
    ll y_max
        = (a * n + b) / m, x_max = (y_max * m - b);
    if (y_max == 0) return ans;
    ans += (n - (x_max + a - 1) / a) * y_max;
    ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
    return ans;
}// sum^{{
        n-1}_0 floor((a * i + b) / m) in log(n + m + a + b)
```

6.10 Miller Rabin* [06308c]

```
while (--t)
     if ((x = mul(x, x, n)) = n - 1) return 1;
  return 0:
}
6.11 Big number [6d475b]
template < typename T >
inline string to_string(const T& x){
  stringstream ss;
  return ss << x, ss. str();
struct bigN:vector<ll>{
  const static int base=1000000000, width=log10(base);
  bool negative;
  bigN(const\_iterator
         a, const_iterator b): vector < ll > (a, b) {}
  bigN(string s){
     if(s.empty())return;
if(s[0]=='-')negative=1,s=s.substr(1);
     else negative=0;
     \label{eq:formula} \text{for} \, (\, \text{int} \, \, \, i \! = \! \text{int} \, (\, s \, . \, \, \text{size} \, (\,) \,) \, \text{-} \, 1 \, ; i \! > \! = \! 0 ; i \! - \! = \! \text{width} \, ) \, \{
       l\hat{l} t=0;
       for(int j=max(0,i-width+1);j \le i;++j)
          t=t*10+s[j]-'0';
       push_back(t);
     trim();
  template<typename T>
     bigN(const T &x):bigN(to_string(x)){}
  bigN():negative(0)\{\}
   void trim(){
     while(size()&&!back())pop_back();
     if (empty()) negative=0;
   void carry(int _base=base){
     for(size_t i=0; i < size(); ++i){
       if (at (i)>=0&&at (i)<_base) continue;
        if(i+lu=size())push_back(0);
       int r=at(i)%_base;
       if(r<0)r+=\_base;
       at(i+1)+=(at(i)-r)/\_base, at(i)=r;
   int abscmp(const bigN &b)const{
     if(size()>b.size())return 1;
     if(size()<b.size())return -1;</pre>
     for (int i=int(size())-1;i>=0;--i){
       if (at(i)>b[i]) return 1;
        if (at (i) <b[i]) return -1;
     return 0;
   int cmp(const bigN &b)const{
     if (negative!=b.negative)return negative?-1:1;
     return negative?-abscmp(b):abscmp(b);
   bool operator < (const bigN&b) const {return cmp(b) < 0;}
   bool operator > (const bigN&b) const {return cmp(b) > 0;}
  bool operator <= (const bigN&b) const{return cmp(b) <= 0;}
bool operator >= (const bigN&b) const{return cmp(b) >= 0;}
   bool operator == (const bigN&b) const {return !cmp(b);}
   bool operator!=(const bigN&b)const{return cmp(b)!=0;}
   bigN abs() const {
     bigN res=*this;
     return res.negative=0, res;
  bigN operator - () const {
     bigN res=*this;
     return res.negative=!negative, res.trim(), res;
  bigN operator+(const bigN &b)const{
     if(negative)return -(-(*this)+(-b));
if(b.negative)return *this-(-b);
     bigN res=*this;
     if(b.size()>size())res.resize(b.size());
     for (size_t i=0; i < b \cdot size(); ++i) res[i] +=b[i];
     return res.carry(), res.trim(), res;
  bigN operator-(const bigN &b)const{
     if(negative)return -(-(*this)-(-b));
if(b.negative)return *this+(-b);
      if(abscmp(b)<0)return -(b-(*this));
     bigN res=*this;
     if(b.size()>size())res.resize(b.size());
```

for $(size_t i=0; i< b. size();++i)res[i]-=b[i];$

```
National Yang Ming Chiao Tung University FubukiMyWife
    return res.carry(), res.trim(), res;
  bigN operator*(const bigN &b)const{
    res.negative=negative!=b.negative;
    res.resize(size()+b.size())
    for (size_t i=0; i < size(); ++i)
       for (size_t j=0; j< b. size(); ++j)
         if ((res[i+j]+=at(i)*b[j])>=base){
           res[i+j+1]+=res[i+j]/base;
           res[i+j]%=base;
                                                                     }
        \}//\frac{a}{4}k¥ carry · | · .
    return res.trim(), res;
  bigN operator/(const bigN &b)const{
    int norm=base/(b.back()+1);
    bigN x=abs()*norm;
    bigN y=b.abs()*norm;
    bigN q,r;
    q.resize(x.size());
    for (int i=int(x.size())-1;i>=0;--i){
      r=r*base+x[i];
       int s1=r.size()<=y.size()?0:r[y.size()];</pre>
                                                              | };
       int s2=r.size()<y.size()?0:r[y.size()-1];</pre>
       int d=(ll(base)*s1+s2)/y.back();
       r=r - v*\dot{d};
       while (r.negative)r=r+y, --d;
      q[i]=d;
    q.negative=negative!=b.negative;
    return q.trim(),q;
  bigN operator%(const bigN &b)const{
    return *this-(*this/b)*b;
  friend istream& operator>>(istream &ss, bigN &b) {
    string s;
    return ss \gg s, b=s, ss;
  friend
        ostream& operator << (ostream &ss, const bigN &b) {
    if (b. negative) ss << '-'
    ss << (b.empty()?0:b.back());
    for (int i=int(b.size())-2;i>=0;--i)
                                                               }
      ss<<setw(width)<<setfill('0')<<b[i];
    return ss;
                                                               6.15
  template<typename T>
    operator T(){
      stringstream ss;
      ss << *this;
      T res;
       return ss>>res,res;
};
6.12 Fraction [4ab37a]
struct fraction {
                                                                 <del>||m</del>;
  ll n, d;
  fraction
       (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
    11 t = \gcd(n, d);
    n /= t, d /= t;
    if'(d < 0)'n = -n, d = -d;
  fraction operator - () const
  { return fraction(-n, d);
  fraction operator+(const fraction &b) const
  { return fraction(n * b.d + b.n * d, d * b.d); }
  fraction operator-(const fraction &b) const
  { return fraction(n * b.d - b.n * d, d * b.d); } fraction operator*(const fraction &b) const
  { return fraction(n * b.n, d * b.d); }
  fraction operator/(const fraction &b) const
  { return fraction(n * b.d, d * b.n); }
  void print() {
    cout << n;
    if (d != 1) cout << "/" << d;
  }
                                                                     }
6.13 Simultaneous Equations [a231be]
struct matrix { //m variables, n equations
  int n, m;
```

 $fraction\ M[MAXN] [MAXN + 1],\ sol [MAXN];$

int solve() { // -1: inconsistent , >= 0: rank for (int i = 0; i < n; ++i) { int piv = 0; while (piv < m && !M[i][piv].n) ++piv; if (piv == m) continue; for (int j = 0; j < n; ++j) { if (i == j) continue; $\begin{array}{ll} fraction \ tmp = -M[\,j\,] \big[\,piv\,] \ / \ M[\,i\,] \big[\,piv\,] \,; \\ for \ (int \ k = 0; \ k <= \end{array}$ m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];int rank = 0;for (int i = 0; i < n; ++i) { int piv = 0; $\label{eq:while} \mbox{ while $(piv < m \&\& !M[i][piv].n) ++piv;}$ if (piv = m & M[i][m].n) return -1; else if (piv < m) ++rank, sol[piv] = M[i][m] / M[i][piv]; return rank; 6.14 Pollard Rho* [fdef9b] map<ll , int> cnt;void PollardRho(ll n) { if (n = 1) return; if (prime(n)) return ++cnt[n], void(); if (n % 2 = 0) return PollardRho(n / 2), ++cnt[2], void(); 11 x = 2, y = 2, d = 1, p = 1;#define f(x, n, p) ((mul(x, x, n) + p) % n) while (true) { if (d != n & d != 1) { PollardRho(n / d); PollardRho(d); return: if (d = n) ++p; $x \, = \, f \, (\, x \, , \, \, n \, , \, \, p) \, , \ \, y \, = \, f \, (\, f \, (\, y \, , \, \, n \, , \, \, p) \, , \, \, n \, , \, \, p) \, ;$ d = gcd(abs(x - y), n);Simplex Algorithm [6b4566] const int MAXN = 11000, MAXM = 405; const double eps = 1E-10; double a [MAXN] [MAXM], b [MAXN], c [MAXM]; double d [MAXN] [MAXM], x [MAXM]; int ix [MAXN+MAXM]; // !!! array all indexed from 0 // max{cx} subject to {Ax=b,x>=0} // n: constraints, m: vars !!! // x[] is the optimal solution vector // usage : // value = simplex(a, b, c, N, M); double simplex(int n, int m){ $fill_n(d[n], m + 1, 0);$ $fill_n(d[n + 1], m + 1, 0);$ iota(ix, ix + n + m, 0);int r = n, s = m - 1;for (int i = 0; i < n; ++i) { for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j]; d[i][m - 1] = 1; d[i][m] = b[i];if (d[r][m] > d[i][m]) r = i; $copy_n(c, m - 1, d[n]);$ d[n + 1][m - 1] = -1;for (double dd;;) { if (r < n)swap(ix[s], ix[r+m]); $\begin{aligned} &\text{d}[r][s] = 1.0 \ / \ d[r][s]; \\ &\text{for (int } j = 0; \ j <= m; \ ++j) \\ &\text{if } (j := s) \ d[r][j] \ ^*= -d[r][s]; \end{aligned}$ for (int i = 0; i <= n + 1; ++i) if (i != r) { for (int j = 0; j <= m; ++j) if (j != s) d[i][j] += d[r][j] * d[i][s]; d[i][s] *= d[r][s];r = s = -1;for (int j = 0; j < m; ++j) if (s < 0 | | ix[s] > ix[j]) {

```
\begin{array}{l} if \ (d[n+1][j] > eps \ | \ | \\ \ (d[n+1][j] > -eps \ \&\& \ d[n][j] > eps)) \end{array}
   if (s < 0) break;
   for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
     if (r < 0 | |
           (dd = d[r][m]
                   \frac{1}{d[r][s]} - d[i][m] / d[i][s] < -eps | |
           (dd < eps \&\& ix[r+m] > ix[i+m]))
        r = i;
   if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not executable
double ans = 0:
fill_n(x, m, 0);
for (int i = m; i <
  \begin{array}{ll} n+m; ++i) & \text{if } (ix[i] = 0) \\ \text{if } (ix[i] < m-1) \\ \text{ans } += d[i-m][m] & \text{c[ix[i]]}; \\ \text{c[ix[i]]} & \text{c[ix[i]]}; \\ \end{array}
     x[ix[i]] = d[i-m][m];
  }
return ans;
```

6.15.1 Construction

| Primal | Dual |
|--|--|
| Maximize $c^{\intercal}x$ s.t. $Ax \leq b, x \geq 0$ | Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \ge c, y \ge 0$ |
| Maximize $c^{\intercal}x$ s.t. $Ax \leq b$ | Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y = c, y \ge 0$ |
| Maximize $c^{T}x$ s.t. $Ax = b, x \ge 0$ | Minimize $b^{T}y$ s.t. $A^{T}y \ge c$ |

 $\overline{\mathbf{x}}$ and $\overline{\mathbf{y}}$ are optimalified only if for all $i \in [1, n]$, either $\overline{x}_i = 0$ or $\sum_{j=1}^m A_{ji} \overline{y}_j = c_i$ holds and for all $i \in [1, m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij} \bar{x}_j = b_j$ holds.

- 1. In case of minimization, let $c_i' = -c_i$
- $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$ 2.
- 3. $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ $\bullet \sum_{1 \le i \le n} A_{ji} x_i \le b_j$ $\bullet \sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.16 chineseRemainder [a53b6d]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
  ll g = gcd(m1, m2);
if ((x2 - x1) \% g) return -1; // no sol
  m1 /= g; m2 /= g;
  pll p = exgcd(m1, m2);
  ll lcm = m1 * m2 * g;

ll res = p.first * (x^2 - x^1) * m1 + x1;
  // be careful with overflow
  return (res % lcm + lcm) % lcm;
```

6.17 Factorial without prime factor* [c324f3]

```
// O(p^k + log^2 n), pk = p^k
ll prod [MAXP]
ll fac_no_p(ll n, ll p, ll pk) {
   \operatorname{prod}[0] = 1;
   for (int i = 1; i \le pk; ++i)
     if (i % p) prod[i] = prod[i - 1] * i % pk;
      else prod[i] = prod[i - 1];
   11 \text{ rt} = 1;
  \begin{array}{lll} & for & (; \ n; \ n \ / = \ p) \ \{ \\ & rt \ = \ rt \ * \ mpow(prod [\ pk] \ , \ n \ / \ pk , \ pk) \ \% \ pk; \end{array}
      rt = rt * prod[n % pk] % pk;
} // (n! without factor p) % p^k
```

6.18 PiCount* [cad6d4]

```
ll PrimeCount(ll n) { // n ~ 10^13 \Rightarrow < 2s
  if (n <= 1) return 0;
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector < int > smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i \le v; ++i) smalls[i] = (i + 1) / 2;
  for (int i = 0; i < s; ++i) {
    roughs[i] = 2 * i + 1;
larges[i] = (n / (2 * i + 1) + 1) / 2;
      (int p = 3; p \le v; ++p) {
    if (smalls[p] > smalls[p - 1]) {
  int q = p * p;
      ++pc;
```

```
if (1LL * q * q > n) break;
     skip[p] = 1;
     for (int i = q; i \le v; i += 2 * p) skip [i] = 1;
     int ns = 0;
     for (int k = 0; k < s; ++k) {
       int i = roughs[k];
       if (skip[i]) continue;

ll d = lLL * i * p;

larges[ns] = larges[k] - (d <= v ? larges
            [smalls[d] - pc] : smalls[n / d]) + pc;
       roughs [ns++] = i;
     }
     s = ns;
     for (int j = v / p; j >= p; --j) {
      int c =
       }
for (int k = 1; k < s; ++k) {
  const ll m = n / roughs[k];
  ll t = larges[k] - (pc + k - 1);</pre>
  for (int l = 1; l < k; ++l) {
    int p = roughs[1];

if (1LL * p * p > m) break;

t -= smalls[m / p] - (pc + 1 - 1);
  larges[0] = t;
return larges[0];
```

6.19 Discrete Log* [da27bf]

```
int \ Discrete Log(int \ s, \ int \ x, \ int \ y, \ int \ m) \ \{
   constexpr int kStep = 32000;
  unordered\_map < int, int > p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {
    p[y] = i;

y = 1LL * y * x % m;

b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m = 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {
     if (s == y) return i;
s = 1LL * s * x % m;
   if (s = y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
    if \ (fpow(x, p, m) \mathrel{!=} y) \ return \ \text{-}1; \\
   return p;
```

6.20 Berlekamp Massey [3eb6fa]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
   vector < T > d(SZ(output) + 1), me, he;
   for (int f = 0, i = 1; i <= SZ(output); ++i) {
    for (int j = 0; j < SZ(me); ++j)
        d[i] += output[i - j - 2] * me[j];
    if ((d[i] -= output[i - 1]) == 0) continue;

      if (me.empty()) {
        me. resize(f = i);
        continue:
     return me;
```

6.21 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 \ 98789101 \ 987777733 \ 999991921 \ 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

6.22 Estimation

n | 2345 6 7 8 9 20 30 40 50 100 p(n) | 23571115223062756044e42e52e8n | 1001e31e6 1e9 1e12 1e15 1e18 d(i) 12 32 2401344672026880103680 $n \mid 12345678$ 9 10 11 12 13 14 15 $\binom{2n}{n}$ 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

6.23 General Purpose Numbers

• Bernoullinumbers

$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. kj:ss.t. $\pi(j) > \pi(j+1), k+1j$:ss.t. $\pi(j) \ge j$, kj:ss.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.24 Tips for Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$
 - $-A(rx) \Rightarrow r^n a_n$
 - $-A(x)+B(x) \Rightarrow a_n+b_n$
 - $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $-xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$
 - $-A(x)+B(x) \Rightarrow a_n+b_n$
 - $-A^{(k)}(x) \!\Rightarrow\! a_{n+k}$
 - $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
 - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $-xA(x) \Rightarrow na_n$
- Special Generating Function
 - $(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$

$-\frac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{i}{n-1} x^i$ **Polynomial**

Fast Fourier Transform [5e2ea2]

```
const int \max = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = acos(-1);
cplx omega[maxn + 1];
\begin{array}{c} \text{void prefft() } \{\\ \text{ for (int } i = 0; i <= maxn \end{array}
           ; ++i) \text{ omega}[i] = \exp(i * 2 * pi / maxn * I);
void bin(vector < cplx > &a, int n) {
     int lg;
```

```
for (lg = 0; (1 \ll lg) < n; ++lg); --lg;
    vector < cplx > tmp(n);
    for (int i = 0; i < n; ++i) {
         int to = 0;
         for (int j = 0; (1 << j) <
            n; ++j) to |= (((i >> j) & 1) << (lg - j));
        tmp[to] = a[i];
    for (int i = 0; i < n; ++i) a[i] = tmp[i];
void fft (vector<cplx> &a, int n) {
    bin(a, n);
    for (int step = 2; step \leq n; step \leq 1) {
         int to = step \gg 1;
         for (int i = 0; i < n; i += step) {
             for (int k = 0; k < to; ++k) {
                 cplx x = a[i
                     + \text{ to } + \text{ k}] * omega[maxn / step * k];
                 a[i + to + k] = a[i + k] - x;
                 a[i + k] += x;
        }
    }
}
void ifft (vector<cplx> &a, int n) {
    fft(a, n);
    reverse(a.begin() + 1, a.end());
    for (int i = 0; i < n; ++i) a[i] /= n;
vector<int> multiply(const vector<
    int> &a, const vector<int> &b, bool trim = false) {
    int d = 1;
    while
        (d < max(a.size(), b.size())) d <<= 1; d <<= 1;
    vector < cplx > pa(d), pb(d);
    for (int i
         = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
    for (int i
         = \ 0; \ i < b. \, size(); \ +\!\!\!+\!\! i) \ pb[i] = cplx(b[i], \ 0);
    fft(pa, d); fft(pb, d);
    for (int i = 0; i < d; ++i) pa[i] *= pb[i];
    ifft (pa, d);
    vector < int > r(d);
    for (int
          i = 0; i < d; ++i) r[i] = round(pa[i].real());
    if (trim)
         while (r.size() \& r.back() = 0) r.pop_back();
    return r;
 Prime
             Root
                   {\rm Prime}
                                Root
 7681
             17
                    167772161
                                3
                   104857601
 12289
                                3
             11
                   985661441
 40961
             3
                                3
 65537
            3
                   998244353
                                3
 786433
            10
                    1107296257
                                10
 5767169
            3
                   2013265921
                                31
 7340033
                   2810183681
                                11
 23068673
                   2885681153
            3
                                3
 469762049
                   605028353
                                3
7.2 Number Theory Transform* [7d51db]
vector<int> omega;
void Init() {
  omega. resize(kN + 1);
```

```
long long x = fpow(kRoot, (Mod - 1) / kN);
  omega[0] = 1;

for (int i = 1; i <= kN; ++i) {

  omega[i] = 1LL * omega[i - 1] * x % kMod;
void Transform(vector<int> &v, int n) {
  BitReverse(v, n);
for (int s = 2; s \ll n; s \ll 1) {
    int z = s \gg 1;
     for (int i = 0; i < n; i += s) {
       for (int k = 0; k < z; ++k) {
         int x = 1LL
                * v[i + k + z] * omega[kN / s * k] % kMod;
          v[i + k + z] = (v[i + k] + kMod - x) \% kMod;
         (v[i + k] += x) \% = kMod;
    }
  }
```

7.3 Fast Walsh Transform* [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
for (int L = 2; L <= n; L <<= 1)</pre>
      for (int i = 0; i < n; i += L)
         for (int j = i; j < i + (L >> 1); ++j)

a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f
      N = [1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
void
      subset_convolution(int *a, int *b, int *c, int L) {
    // c_k = \sum_{i = 1}^{n} i \cdot j = k, i \cdot k \cdot j = 0 a_i * b_j
   int n = 1 \ll L;
   for (int i = 1; i < n; ++i)
   ct[i] = ct[i & (i - 1)] + 1;
for (int i = 0; i < n; ++i)
   \begin{array}{l} f[\,ct\,[\,i\,]\,][\,i\,] \,=\, a\,[\,i\,]\,,\,\, g[\,ct\,[\,i\,]\,][\,i\,] \,=\, b\,[\,i\,]\,;\\ for\,\,\, (\,int\,\,i\,=\,0\,;\,\,i\,<=\,L\,;\,\,+\!\!+\!\!i\,) \end{array}
      fwt(f[i], n, 1), fwt(g[i], n, 1);
   for (int i = 0; i \le L; ++i)
      for (int j = 0; j <= i; ++j)
          for (int x = 0; x < n; ++x)

h[i][x] += f[j][x] * g[i - j][x];
   for (int i = 0; i <= L; ++i)

fwt(h[i], n, -1);

for (int i = 0; i < n; ++i)
      c[i] = h[ct[i]][i];
```

7.4 Polynomial Operation [869cb1]

```
\begin{array}{ll} fi(s,\ n) \ for \ (int\ i=(int)(s); \ i<(int)(n); \ +\!\!+\!\!i) \\ template<int\ MAXN, \ ll\ P, \ ll\ RT>//\ MAXN=2^k \end{array}
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<ll>::vector;
  static NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
  Poly& irev()
        \{ return reverse(data(), data() + n()), *this; \}
  Poly& isz(int m) { return resize(m), *this; }
  Poly& iadd(const Poly &rhs) \{ // n() = rhs.n() \}
    fi(0, n()) if
  (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
  Poly& imul(ll k) {
     fi(0, n()) (*this)[i] = (*this)[i] * k % P;
return *this;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
     while (m < n() + rhs.n() - 1) m <<= 1;
     Poly X(*this, m), Y(rhs, m);
     \begin{array}{l} {\rm ntt}\,(X.\,{\rm data}\,()\;,\;m)\;,\;\;{\rm ntt}\,(Y.\,{\rm data}\,()\;,\;m)\;;\\ {\rm fi}\,(\,0\;,\;m)\;\;X[\;i\;]\;=\;X[\;i\;]\;\;^*\;Y[\;i\;]\;\%\;P; \end{array}
     ntt(X.data(), m, true);
     return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
     if (n() = 1) return \{ntt.minv((*this)[0])\};
     int m = 1;
     while (m < n() * 2) m <<= 1;
     Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
     Poly Y(*this, m);
     ntt(Xi.data(), m), ntt(Y.data(), m);
     fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
       if ((Xi[i]) = P) < 0) Xi[i] += P;
```

```
ntt(Xi.data(), m, true);
  return Xi. isz(n());
Poly Sqrt()
      const { // Jacobi((*this)[0], P) = 1, 1e5/235ms}
  if (n()
       = 1) return {QuadraticResidue((*this)[0], P)};
  Poly
       X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
  return
        X.\,iadd\left(Mul(X.\,Inv\left(\right)\right).\,is\,z\left(n\left(\right)\right)\right).\,imul\left(P\ /\ 2\ +\ 1\right);
pair < Poly , Poly > DivMod
  (const Poly &rhs) const { // (rhs.)back() != 0 if (n() < rhs.n()) return {{0}, *this}; const int m = n() - rhs.n() + 1;
  Poly X(rhs); X.irev().isz(m);
Poly Y(*this); Y.irev().isz(m);
 Poly Q = Y.Mul(X.Inv()).isz(m).irev();

X = \text{rhs.Mul}(Q), Y = *\text{this};

fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0,
       ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n())
        ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
  return ret;
Poly
       _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn -
  return Poly(\dot{Y}.data() + \dot{n}() - 1, \dot{Y}.data() + \dot{Y}.\dot{n}());
vector<ll> _eval(const
      vector<ll> &x, const vector<Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return { };
  vector Poly> down (m * 2);
  // down[1] = DivMod(up[1]) . second;
// fi(2, m *
        2) down[i] = down[i / 2].DivMod(up[i]).second;
  down[1] = Poly(up[1])
  1]._tmul(up[i].n() - 1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
vector<Poly> up(m * 2);
  fi(0, m) up[m+i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i
       > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
vector
    <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate (const vector
    <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
fi(0, m) down[m + i] = {z[i]};
  for (int i = m -
       `1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
  return Dx(). Mul(Inv()). Sx(). isz(n());
Poly Exp() const { // (*this)[0] = 0, 1e5/360ms
  if (n() = 1) return \{1\};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
```

```
Poly Y = X.Ln(); Y[0] = P - 1;
    \begin{array}{c} fi\,(\,0\,,\,\,n\,(\,)\,) \\ if\,\,\,(\,(Y[\,i\,]\,=\,(\,^*t\,his\,)\,[\,i\,]\,\,\,-\,\,Y[\,i\,]\,)\,<\,\,0)\,\,Y[\,i\,]\,\,+\!\!=\,P; \end{array}
    return X.Mul(Y).isz(n());
   / M := P(P - 1). If k >= M, k := k \% M + M.
  Poly Pow(ll k) const {
    int nz = 0;
    while (nz < n() && !(*this)[nz]) ++nz;
    if (nz * min(k, (ll)n()) >= n()) return Poly(n());
    if (!k) return Poly(Poly {1}, n());
Poly X(data() + nz, data() + nz + n() - nz * k);
    return X.Ln().imul
         (k % P).Exp().imul(c).irev().isz(n()).irev();
  static 11
       LinearRecursion(const vector<ll> &a, const vector
       \langle ll \rangle \& coef, ll n \rangle \{ // a_n = \langle sum c_j a_(n-j) \rangle \}
    const int k = (int)a.size();
    assert((int)coef.size() = k + 1);
    while (n) {
      if (n \% 2) W = W.Mul(M).DivMod(C).second;
      n \neq 2, M = M.Mul(M).DivMod(C).second;
    fi(0, k) ret = (ret + W[i] * a[i]) \% P;
    return ret;
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template \Leftrightarrow \ decltype(Poly\_t::ntt) \ Poly\_t::ntt = \{\,\};
```

7.5 Value Polynomial [96cde9]

```
struct Poly {
  mint base; // f(x) = poly[x - base]
  vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
  mint get_val(const mint &x) {
    if (x >= base \&\& x < base + SZ(poly))
       return poly[x - base];
    mint rt = 0;
     vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
    for (int i = 1; i < SZ(poly); ++i)
    for (int i = 0; i < SZ(poly); ++i)
rt += poly[i] * ifac[i] * inegfac
[SZ(poly) - 1 - i] * lmul[i] * rmul[i];
  void raise() { // g(x) = sigma{base:x} f(x)
    if (SZ(poly) == 1 && poly[0] == 0)
       return;
    mint nw = get\_val(base + SZ(poly));
    poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)
       poly[i] += poly[i - 1];
};
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P) = 0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

8 Geometry

8.1 Basic [e38806]

```
bool same
    (double a, double b) { return abs(a - b) < eps; }
struct P {</pre>
```

```
double x, y;
   P() : x(0), y(0) \{ \}
  P(): x(0), y(0) {}
P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }
double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
double abs() { return bypost(x y); }
   double abs() { return hypot(x, y); }
   P unit() { return *this / abs(); }
   P rot(double o) {
      double c = cos(o), s = sin(o);
return P(c * x - s * y, s * x + c * y);
   double angle() { return atan2(y, x); }
struct L {
   // ax + by + c = 0
   double a, b, c, o;
   P pa, pb;
   L() : a(0), b(0), c(0), o(0), pa(), pb() \{ \}
   L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x ), c(pa pb), o(atan2(-a, b)), pa(pa), pb(pb) {}
     project (P p) { return pa + (pb - pa) . unit
() * ((pb - pa) * (p - pa) / (pb - pa) . abs()); }
  () ((pb - pa) (p - pa) / (pb - pa).abs()), }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs() * (pb - pa).abs()); }
   bool inside (
         P p) \{ return min(pa.x, pb.x) \le p.x \& p.x \le max \}
         (pa.x, pb.x) && min(pa.y, pb.y) \le p.y && p.y \le max(pa.y, pb.y) && same(a * p.x + b * p.y, -c);}
};
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
   if (\max(p1.x, p2.x) < \min(p3.x, p4.x) | |
          \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
   \inf (\max(p1.y, p2.y) < \min(p3.y, p4.y) | |
          \max(\texttt{p3.y}, \texttt{p4.y}) < \min(\texttt{p1.y}, \texttt{p2.y})) \ \textbf{return} \ \textbf{false} \,;
   return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^ (p4 - p2)) <= 0 && sign((p1 - p3) ^
              (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
      (L x, L y) { return same(x.a * y.b, x.b * y.a); }
P Intersect
      (L x, L y) { return P(-x.b * y.c + x.c * y.b, x .a * y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
8.2 KD Tree [375ca2]
namespace kdt {
int root, lc [maxn],
        rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
   if (1 = r) return -1;
   function {<} bool(const\ point\ \&,\ const\ point
          &>> f = [dep](const point &a, const point &b) {
       if (dep \& 1) return a.x < b.x;
      else return a.y < b.y;
   int m = (l + r) \gg 1;

nth\_element(p + l, p + m, p + r, f);
   xl[m] = xr[m] = p[m].x;
   yl[m] = yr[m] = p[m].y;
   lc[m] = build(1, m, dep + 1);
   if (~lc[m]) {
      xl[m] = min(xl[m], xl[lc[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
   rc[m] = build(m + 1, r, dep + 1);
   if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
   return m;
```

```
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds | | q.x > xr[o] + ds | |
       q.y <
            yl[o] - ds \mid \mid q.y > yr[o] + ds) return false;
  return true:
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
void dfs (
     const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
if (cd != 0) d = min(d, cd);
  if ((dep & 1)
     && q.x < p[o].x \mid | !(dep & 1) && q.y < p[o].y) { if <math>(\sim lc[o]) dfs(q, d, lc[o], dep + 1); if (\sim lc[o]) dfs(q, d, rc[o], dep + 1);
     if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
     if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
     if (\sim lc [o]) dfs(q, d, lc [o], dep + 1);
void init (const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
long long nearest (const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
     Sector Area [c41fb7]
  calc area of sector which include a, b
```

```
// calc area of sector which include a, b double SectorArea(P a, P b, double r) { double o = atan2(a.y, a.x) - atan2(b.y, b.x); while (o <= 0) o += 2 * pi; while (o >= 2 * pi) o -= 2 * pi; o = min(o, 2 * pi - o); return r * r * o / 2; }
```

8.4 Half Plane Intersection [f7274e]

```
bool jizz (L l1, L l2, L l3) {
  P = Intersect(12, 13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;
bool cmp(const L &a, const L &b){
  return same(
       a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls) {
  sort(ls.begin(),ls.end(),cmp);
  vector < L > pls(1, ls[0]);
  for (int i=0; i<(int) ls. size ();++i) if (!
        same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
  deque < int > dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c
     ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
  for (int i=2; i < (int) pls. size(); ++i){
    \begin{array}{l} meow(\,i\,\,,dq\,.\,back\,(\,)\,\,,dq\,[\,dq\,.\,size\,(\,)\,\,-\,2\,]\,)dq\,.\,pop\_back\,(\,)\,\,;\\ meow(\,i\,\,,dq\,[\,0\,]\,,dq\,[\,1\,]\,)dq\,.\,pop\_front\,(\,)\,\,; \end{array}
     dq.push_back(i);
  meow(dq
        .front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop\_front();
  if (dq.size()<3u)return vector
        <P>(); // no solution or solution is not a convex
  vector<P> rt;
  for (int i=0; i<(int)dq. size();++i)rt.push_back
        (Intersect (pls [dq[i]], pls [dq[(i+1)%dq.size()]]));
}
```

8.5 Rotating Sweep Line [0411f0]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
  int n=int(ps.size());
```

```
vector < int > id(n), pos(n);
  vector{<}pair{<}int\;,int{>\!>} line\left(n^*(n\text{-}1)/2\right);
   int m=-1
   for(int i=0;i<n;++i)for
       (int j=i+1;j< n;++j)line[++m]=make\_pair(i,j); ++m;
   sort(line.begin(),line.end(),[&](const
         \verb"pair<int", \verb"int> \&a", \verb"const" pair<int", \verb"int> \&b")->bool\{
       if (ps
            [a.first].first=ps[a.second].first)return 0;
       if (ps
            [b.first].first = ps[b.second].first)return 1;
       return (double
            ) (ps[a.first].second-ps[a.second].second)/(ps
             [a.first].first-ps[a.second].first) < (double
            )(ps[b.first].second-ps[b.second].second
            )/(ps[b.first].first-ps[b.second].first);
       });
  for (int i=0;i< n;++i)id[i]=i;
   sort(id.begin(),id.end(),[&](const
        int &a, const int &b){ return ps[a]<ps[b]; });
   for (int i=0; i< n; ++i) pos [id[i]]=i;
  for (int i=0; i < m++i){
     auto l=line[i];
     // meow
     \texttt{tie} \, (\, \texttt{pos} \, [\, \texttt{l.first} \, ] \, , \texttt{pos} \, [\, \texttt{l.second} \, ] \, ,
          id [pos[1.first]], id [pos[1.second]])=make_tuple
          (pos[l.second],pos[l.first],l.second,l.first);
}
```

8.6 Triangle Center [4e8ee9]

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) /
  double ay = (a.y + b.y)
  double bx = (c.x + b.x) /
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
 return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b
      , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);
     la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (
la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
```

8.7 Polygon Center [ee6ff0]

```
Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

8.8 Maximum Triangle [3a6d38]

```
double ConvexHullMaxTriangleArea
    (Point p[] , int res[] , int chnum) {
```

-eps) $t.emplace_back(a.x + i * x, a.y + i * y);$

```
if (j - 1.0 \le eps \& j > =
   double area = 0, tmp:
                                                                                     -eps) t.emplace_back(a.x + j * x, a.y + j * y);
   res[chnum] = res[0];
   for (int i = 0, j = 1, k = 2; i < chnum; i++) {
     while (fabs(Cross(p[
                                                                            return t;
          res[j] - p[res[i]] , p[res[(k+1) \% chnum]]
          \begin{array}{l} p[\operatorname{res}[i]])) > \operatorname{fabs}(\operatorname{Cross}(p[\operatorname{res}[j]] - p[\operatorname{res}[i]], \\ p[\operatorname{res}[k]] - p[\operatorname{res}[i]]))) \ k = (k+1) \% \ \text{chnum}; \end{array}
                                                                         // calc area
                                                                                intersect by circle with radius \boldsymbol{r} and triangle O\!AB
                                                                          double AreaOfCircleTriangle(P a, P b, double r) {
     tmp = fabs(Cross(
          p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                                                            bool ina = a.abs() < r, inb = b.abs() < r;
     if (tmp > area) area = tmp;
                                                                            auto p = CircleCrossLine(a, b, P(0, 0), r);
     while (fabs (Cross (p[
                                                                            if (ina) {
           \begin{array}{l} {\rm res}\,[(\,j+1)\,\%\,\,{\rm chnum}\,] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,\, , \,\, p[\,{\rm res}\,[\,k\,]] \,\, - \,\, \\ p[\,{\rm res}\,[\,i\,]])) > \,\, {\rm fabs}\,(\,{\rm Cross}\,(p[\,{\rm res}\,[\,j\,]] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,\, , \\ p[\,{\rm res}\,[\,k\,]] \,\, - \,\, p[\,{\rm res}\,[\,i\,]]))) \,\,\, j \,\, = \,\, (j\,+\,1)\,\,\%\,\,{\rm chnum}; \\ \end{array} 
                                                                               if (inb) return abs(a ^ b) / 2;
                                                                               return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
     tmp = fabs (Cross (
                                                                             if (inb) return
                                                                                  SectorArea(p[0], a, r) + abs(p[0] ^ b)
          p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                                                            if (p.size() = 2u) return SectorArea(a, p[0], + SectorArea(p[1], b, r) + abs(p[0] p[1]
     if (tmp > area) area = tmp;
                                                                                                                                  p[1]) / 2;
                                                                             else return SectorArea(a, b, r);
  return area / 2;
                                                                          // for any triangle
8.9 Point in Polygon [0a9a66]
                                                                         double AreaOfCircleTriangle(vector<P> ps, double r) {
                                                                            double ans = 0;
int pip(vector<P> ps, P p) {
                                                                            for (int i = 0; i < 3; ++i) {
int j = (i + 1) \% 3;
   for (int i = 0; i < ps.size(); ++i) {
                                                                               double o = atan2
     int a = i, b = (i + 1) \% ps.size();
                                                                               L l(ps[a], ps[b]);
     P q = l.project(p);
     if ((p - q).abs() < eps && l.inside(q)) return 1;
                                                                               ans += AreaOfCircleTriangle
     if (same(ps[
                                                                                    (\,ps\,[\,i\,]\;,\;\;ps\,[\,j\,]\;,\;\;r\,)\;\;*\;\;(\,o>=\;0\;\;?\;\;1\;\;:\;\;-1)\,;
          a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
     if (ps[a].y > ps[b].y) swap(a, b);
if (ps[a].y \le p.y && p.y <
                                                                            return abs(ans);
          ps[b].y \& p.x \le ps[a].x + (ps[b].x - ps[a].x 
) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) ++c;
                                                                          8.11
                                                                                   Tangent of Circles and Points to Circle
                                                                                   [19eb58]
   return (c & 1) * 2;
                                                                          vector <\!\!L\!\!> tangent (C~a,~C~b)~\{
                                                                         #define Pij \
8.10 Circle [466c44]
                                                                            P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); \
                                                                            z.emplace\_back(a.c + i, a.c + i + j);
struct C {
                                                                          #define deo(I,J)
  Р с;
                                                                            double d = (a)
   double r
  C(P \ c = P(0, 0), \ double \ r = 0) : c(c), \ r(r) \ \{\}
                                                                                 .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
                                                                            vector<P> Intersect(C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
                                                                            \quad \text{if} \ (a.r < b.r) \ swap(a, \ b); \\
                                                                            vector < L > z;
   vector<P> p;
                                                                            if ((a.c - b.c).abs() + b.r < a.r) return z;
   if (same(a.r + b.r,
                                                                            else if (same((a.c - b.c).abs() + b.r, a.r)) \{ Pij; \}
        d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
   else if (a.r + b.r > d \&\& d + a.r >= b.r) {
                                                                            else {
     double o = acos
                                                                               deo(-,+);
                                                                               if (same(d, a.r + b.r)) \{ Pij; \}
          ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
     P i = (b.c - a.c).unit();
                                                                               else if (d > a.r + b.r) \{ deo(+,-); \}
     p.push\_back(a.c + i.rot(o) * a.r);
     p.push\_back(a.c + i.rot(-o) * a.r);
                                                                            return z;
   return p:
                                                                          vector <L> tangent (C c, P p) {
                                                                            vector<L> z;
double d = (p - c.c).abs();
double IntersectArea(C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
                                                                             if (same(d, c.r)) {
  if (d \ge a.r + b.r - eps) return 0;
if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
                                                                            P i = (p - c.c).rot(pi / 2);
z.emplace_back(p, p + i);
} else if (d > c.r) {
   double p = acos
                                                                               double o = acos(c.r / d);
        ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
   double q = acos
                                                                               P i = (p - c.c).unit
                                                                                    ()\;,\;j\;=\;i\;.\,rot\,(o)\;\;*\;\;c\,.\,r\,,\;\;k\;=\;i\,.\,rot\,(-o)\;\;*\;\;c\,.\,r\,;
  ((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
                                                                               z.\,emplace\_back(\,c.\,c\,\,+\,\,j\,\,,\,\,\,p)\,;
                                                                               z.emplace\_back(c.c + k, p);
// remove second
      level if to get points for line (defalut: segment)
                                                                            return z;
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
                                                                          8.12 Area of Union of Circles [490636]
                                                                          vector<pair<double, double >>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
   double C = sq(a.x - o.x)
       + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
   vector<P> t;
                                                                            vector<pair<double, double>>> res;
if (same(a.r + b.r, d));
   if~(d>=-eps)~\{\\
                                                                            d = \max(0., d);
     double i = (-B - \operatorname{sqrt}(d)) / (2 * A);
     double j = (-B + \operatorname{sqrt}(d)) / (2 * A);
                                                                            else if (d < abs(a.r + b.r) - eps) 
                                                                               if (i - 1.0 <= eps && i >=
```

```
if (z < 0) z += 2 * pi;
     double 1 = z - o, r = z + o;
if (1 < 0) 1 += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
     if (l > r) res.emplace_back
    (l, 2 * pi), res.emplace_back(0, r);
     else res.emplace_back(l, r);
  return res;
double CircleUnionArea
     (vector <C> c) { // circle should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {

vector<pair<double, double>>> s = \{\{2 * pi, 9\}\}\}, z;

for (int j = 0; j < n; ++j) if (i != j) {
        z = CoverSegment(c[i], c[j]);
        for (auto &e : z) s.push_back(e);
     sort(s.begin(), s.end());
auto F = [&] (double t) { return c[i].r * (c[i].r *
            t + c[i].c.x * sin(t) - c[i].c.y * cos(t)); };
     for (auto &e : s) {
        if (e. first > w) a += F(e. first) - F(w);
        w = max(w, e.second);
  return a * 0.5;
```

8.13 Minimun Distance of 2 Polygons [7eb8bb]

8.14 2D Convex Hull [65eaab]

```
bool operator < (const P &a, const P &b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator>(const P &a, const P &b) {
  return same(a.x, b.x) ? a.y > b.y : a.x > b.x;
#define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
 vector<P> p;
 sort(ps.begin(), ps.end(), [&] (P a, P b) { return
    same(a.x, b.x) ? a.y < b.y : a.x < b.x; });</pre>
  for (int i = 0; i < ps.size(); ++i) {
    while (p.size() >= 2 \&\& crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
     (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push\_back(ps[i]);
 p.pop_back();
  return p;
```

```
int sgn(double
      x) \ \{ \ return \ same(x, \ 0) \ ? \ 0 \ : \ x > 0 \ ? \ 1 \ : \ -1; \ \}
P isLL(P p1, P p2, P q1, P q2) {
  double a = \operatorname{crx}(q1, q2, p1), b = \operatorname{-crx}(q1, q2, p2); return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  \text{vector} <\!\!P\!\!> p,\ u,\ d;
  CH() {}
  CH(vector < P > ps) : p(ps) {
    n = ps.size();
     rotate (p. begin
         (), min_element(p.begin(), p.end()), p.end());
     auto t = max\_element(p.begin(), p.end());
     d = vector <\!\!P\!\!>\!\! (p.begin(), next(t));
     u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find (vector <P> &v, P d) {
    int l = 0, r = v.size();
     while (1 + 5 < r) {

int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;

if (v[L] * d > v[R] * d) r = R;
       else \dot{l} = L;
     int x = 1;
     for (int i = l +
           1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
     return x;
  int findFarest(P v) {
     if (v.y > 0 | v.y = 0 \& v.x > 0) return
          ((int)d.size() - 1 + find(u, v)) \% p.size();
     return find(d, v);
    get(int 1, int r, Pa, Pb) {
     int s = sgn(crx(a, b, p[1 \% n]));
     while (1 + 1 < r) {
       else r = m;
     vector <P> getLineIntersect (Pa, Pb) {
    int X = findFarest((b - a).rot(pi / 2));
int Y = findFarest((a - b).rot(pi / 2));
if (X > Y) swap(X, Y);
     if (sgn
           \begin{array}{l} \text{(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0)} \\ \text{return } \{ \text{get(X, Y, a, b), get(Y, X + n, a, b)} \}; \end{array} 
    return {}; // tangent case falls here
  void update_tangent(P q, int i, int &a, int &b) { if (sgn(crx(q, p[a], p[i])) > 0) a = i; if (sgn(crx(q, p[b], p[i])) < 0) b = i;
  void bs(int l, int r, Pq, int &a, int &b) {
     if (l = r) return;
     update_tangent(q, 1 % n, a, b);
      \inf s = sgn(crx(q, p[1 \% n], p[(1 + 1) \% n])); 
     while (l + 1 < r)
       int m = (l + r) \gg 1;
       if (sgn(crx
            (q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
       else r = m;
    update\_tangent(q, r \% n, a, b);
  int x = 1;
  for (int i = l)
        + 1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
  return x;
int findFarest(P v) {
  if (v.y > 0 | | v.y = 0 & v.x > 0) return
         ((int)d.size() - 1 + find(u, v)) \% p.size();
  return find(d, v);
P get(int 1, int r, Pa, Pb) {
  int s = sgn(crx(a, b, p[1 \% n]));
  while (l + 1 < r) {
int m = (l + r) >> 1;
     if (sgn(crx(a, b, p[m\%n])) == s) l = m;
```

```
else r = m:
   return isLL(a, b, p[1 \% n], p[(1 + 1) \% n]);
vector <P> getIS (P a, P b) {
   \begin{array}{ll} \text{int } X = \operatorname{findFarest}((b - a).\operatorname{spin}(pi \ / \ 2));\\ \text{int } Y = \operatorname{findFarest}((a - b).\operatorname{spin}(pi \ / \ 2)); \end{array} 
   if (X > Y) swap(X, Y);
   if \ (\operatorname{sgn}(\operatorname{crx}(a, b, p[X])) * \operatorname{sgn}(\operatorname{crx}(a, b, p[Y])) <
         0) return \{ get(X, Y, a, b), get(Y, X + n, a, b) \};
  return { };
void update_tangent(P q, int i, int &a, int &b) {
  \begin{array}{ll} \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{a}],\mbox{ }p[\mbox{i}])) > 0) \mbox{ }a = \mbox{i}\,;\\ \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{b}],\mbox{ }p[\mbox{i}])) < 0) \mbox{ }b = \mbox{i}\,; \end{array}
void bs(int 1, int r, Pq, int &a, int &b) {
  if (l = r) return;
   update_tangent(q, 1 % n, a, b);
  int s = sgn(crx(q, p[1 \% n], p[(1 + 1) \% n]));
while (1 + 1 < r) {
      int m = (l + r) \gg 1;
      if (sgn
            (crx(q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
      else r = m;
  }
   update\_tangent(q, \ r \ \% \ n, \ a, \ b);
bool contain (P p) {
  if (p.x < d[0].x | | p.x > d.back().x) return 0;
   auto it
        = lower\_bound(d.begin(), d.end(), P(p.x, -1e12));
   \begin{array}{ll} \textbf{if} & (\hspace{1pt} \textbf{it} \hspace{-2pt} \textbf{->} \textbf{x} = \hspace{-2pt} \textbf{p.x}) \end{array} \{
    if (it->y > p.y) return 0;
else if (crx(*prev(it), *it, p) < -eps) return 0;
   i\,t \,=\, lower\_bound
         (u.begin(), u.end(), P(p.x, 1e12), greater < P > ());
   if (it->x = p.x) {
  if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
   return 1;
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
  a = b = 0;
  int i
         = lower\_bound(d.begin(), d.end(), p) - d.begin();
  bs(0, i, p, a, b);
  bs(i, d.size(), p, a, b);
i = lower_bound(
         u.begin(), u.end(), p, greater<P>()) - u.begin();
  bs((int
         d.size() - 1, (int)d.size() - 1 + i, p, a, b);
  bs((int)d.size()
          -1 + i, (int)d.size() -1 + u.size(), p, a, b);
   return 1;
```

8.15 3D Convex Hull [29e4a9]

```
double
       absvol(const P a, const P b, const P c, const P d) {
  return abs(((b-a)^(c-a))*(d-a))/6;
struct convex3D {
  static const int maxn=1010;
  struct T{
     int a,b,c;
     bool res;
     T()\{\}
     T(int a, int
            b, int c, bool res=1: a(a), b(b), c(c), res(res){}
  int n,m;
 P p[maxn];
T f[maxn*8];
  int id [maxn] [maxn]:
  bool on (T &t, P &q) {
     return ((
          p\,[\,t\,.\,c\,]\,\,\hbox{-p}\,[\,t\,.\,b\,]\,)\,\,\widehat{}\,\,(p\,[\,t\,.\,a\,]\,\,\hbox{-p}\,[\,t\,.\,b\,]\,)\,)\,*(q\,\hbox{-p}\,[\,t\,.\,a\,]\,) {>} \mathrm{eps}\,;
  void meow(int q,int a,int b){
     int g=id[a][b];
     if (f [g].res){
        if (on(f[g],p[q]))dfs(q,g);
        else{
```

```
id\left[\,q\,\right]\left[\,b\right]\!=\!id\left[\,a\,\right]\left[\,q\right]\!=\!id\left[\,b\,\right]\left[\,a\right]\!=\!m;
             f[m++]=T(b,a,q,1);
         }
      }
   }
   void dfs(int p,int i){
      f[i].res=0;
       \begin{array}{l} \text{meow(p,f[i].b,f[i].a);} \\ \text{meow(p,f[i].c,f[i].b);} \end{array} 
      meow(p, f[i].a, f[i].c);
    void operator()(){
       if (n<4) return;
       if ([&](){
             for (int i=1; i< n; ++i) if (abs
                   (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
             return
             }() || [&](){
             for (int
                        i=2; i < n; ++i ) if (abs((p[0]-p[i])
                    (p[1]-p[i]) > eps) return swap(p[2],p[i]), 0;
            return 1;
}() || [&](){
             for (int i
                   =3; i < n; ++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                   *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
             return 1;
             }())return;
       for(int i=0;i<4;++i){
         T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
          if(on(t,p[i]))swap(t.b,t.c);
          id [t.a][t.b]=id [t.b][t.c]=id [t.c][t.a]=m;
          f[m++]=t;
       for (int i=4; i< n; ++i) for
             (int j=0; j < m++j) if (f[j].res & on(f[j],p[i])){
          dfs\left( \,i\,\,,\,j\,\right) ;
          break:
       int mm=m; m=0;
      for (int i=0; i < mm + i) if (f[i].res) f[m++]=f[i];
   bool same(int i, int j){
      \textcolor{return}{\textbf{return}} \ ! (\, absvol(\, p\, [\, f\, [\, i\, ]\, .\, a\, ]\, , p\, [\, f\, [\, i\, ]\, .\, a\, ]\, , \\
            ] b) p[f[i].c],p[f[j].a])>eps || absvol(p[f[i].a]),p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].c],p[f[i].c],p[f[j].c])>eps);
   int faces(){
      int r=0;
       for(int i=0;i<m,++i){
          int iden=1;
          for (int j=0; j < i; ++j) if (same(i,j)) iden=0;
          r += iden;
      return r;
} tb;
```

8.16 Minimum Enclosing Circle [fc0e72]

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt \ p0 = b \ - a \, , \ p1 = c \ \cdot
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0
                    p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
}
circle min_enclosing(vector<pt> &p) {
  random\_shuffle(p.begin(),\ p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {
    if (norm2(cent - p[i]) <= r) continue;</pre>
     cent = p[i];
     r = 0.0:
     for (int j = 0; j < i; +++j) {
       if (\text{norm2}(\text{cent} - p[j]) \le r) continue;

\text{cent} = (p[i] + p[j]) / 2;
       cent = (p[i] + p[j])
       r = norm2(p[j] - cent);
for (int k = 0; k < j; ++k)
         if (norm2(cent - p[k]) <= r) continue;</pre>
         cent = center(p[i], p[j], p[k]);
         r = norm2(p[k] - cent);
```

3) $\operatorname{ord}[\operatorname{in}[u]] = \operatorname{ord}[\operatorname{out}[u]] = u$

4) bitset MAXN> inset

```
return circle(cent, sqrt(r));
                                                                                                                                                                     struct Query {
                                                                                                                                                                          int L, R, LBid, lca;
                                                                                                                                                                          Query(int u, int v) {
  int c = LCA(u, v);
                      Closest Pair [f6de57]
 8.17
                                                                                                                                                                                if (c = u \mid | c = v)
 double closest_pair(int 1, int r) {
                                                                                                                                                                                    q.lca \, = \, -1 \, , \, \, q.L = out \, [c \, \, \widehat{} \, \, u \, \, \widehat{} \, \, v] \, , \, \, q.R = out \, [c \, ] \, ;
       // p should be sorted
                                                                                                                                                                                else if (out[u] < in[v])
                      increasingly according to the x-coordinates.
                                                                                                                                                                                     q.lca = c, q.L = out[u], q.R = in[v];
              (l = r) return 1e9;
       if (r - l = 1) return dist(p[l], p[r]);
                                                                                                                                                                               \begin{array}{l} q.\,lca \, = \, c \, , \ q.L \, = \, out \, [\,v\,] \, , \ q.R \, = \, in \, [\,u\,] \, ; \\ q.\,Lid \, = \, q.L \, \, / \  \, blk \, ; \end{array}
       int m = (l + r) >> 1;
       double d =
                     min(closest\_pair(l, m), closest\_pair(m + 1, r));
                                                                                                                                                                          bool operator < (const Query &q) const {
       vector<int> vec;
                                                                                                                                                                                if (LBid != q.LBid) return LBid < q.LBid;
       for (int i = m; i >= 1 &&
                                                                                                                                                                                return R < q.R;
                 \begin{array}{lll} fabs\left(p[m].x - p[i].x\right) < d; & --i\right) & vec.push\_back(i);\\ \left(int & i = m + 1; & i <= r & \& \\ fabs\left(p[m].x\right) & fabs\left(p
       \begin{array}{ll} fabs\left(p\left[m\right].x - p\left[i\right].x\right) < d\,; \, +\!\!\!+\!\!\!i\,\right) \,\, vec.push\_back(\,i\,)\,; \\ sort\left(vec.begin\left(\right), \,\, vec.end\left(\right) \end{array}
                                                                                                                                                                     f,
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
                       [\&](int a, int b) \{ return p[a].y < p[b].y; \});
      for (int i = 0; i < vec.size(); ++i) {
    for (int j = i + 1; j < vec.size()
        && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {
        d = min(d, dist(p[vec[i]], p[vec[j]]));
                                                                                                                                                                                inset[x] = \sim inset[x];
                                                                                                                                                                     void solve(vector<Query> query) {
                                                                                                                                                                          sort(ALL(query));
                                                                                                                                                                          int L = 0, R = 0;
                                                                                                                                                                          for (auto q : query) {
       return d;
                                                                                                                                                                                while (R < q.R) flip (ord[++R]);
                                                                                                                                                                                while (L > q.L) flip (ord[-L]);
                                                                                                                                                                                while (R > q.R) flip (ord[R--]);
              Else
 9
                                                                                                                                                                                \label{eq:while} \text{while } (L < q.L) \ \text{flip} (\text{ord} [L++]);
                  Cyclic Ternary Search* [9017cc]
                                                                                                                                                                                if (~q.lca) add(arr[q.lca]);
                                                                                                                                                                                // answer query
                                                                                                                                                                                if (~q.lca) sub(arr[q.lca]);
    * bool pred(int a, int b);
 f(0) \sim f(n-1) is a cyclic-shift U-function
                                                                                                                                                                  }
  return idx s.t. pred(x, idx) is false for all x*/
  int cyc_tsearch(int n, auto pred) {
                                                                                                                                                                    9.4
                                                                                                                                                                                   Additional Mo's Algorithm Trick
       if (n = 1) return 0;
                                                                                                                                                                          Mo's Algorithm With Addition Only
       int l = 0, r = n; bool rv = pred(1, 0); while (r - l > 1) {

    Sort queryssame as the normal Mo's algorithm.

                                                                                                                                                                                   For each query [l,r]:
            int m = (1 + r) / 2;
                                                                                                                                                                              - If l/blk = r/blk, brute-force.
              \begin{tabular}{ll} \be
                                                                                                                                                                              - If l/blk \neq curL/blk, initialize curL := (l/blk + 1) \cdot blk, curR :=
             else l = m;
                                                                                                                                                                                    curL-1
                                                                                                                                                                              - If r > curR, increase curR
       return pred(1, r % n) ? 1 : r % n;
                                                                                                                                                                                  decrease cur L to fit l, and then undo after answering
                                                                                                                                                                         Mo's Algorithm With Offline Second Time
                                              Algorithm(With
                                                                                                               modification)
 9.2
                  Mo's
                                                                                                                                                                              - Require: Changing answer \equiv adding f([l,r],r+1).
                                                                                                                                                                              - Require: f([l,r],r+1) = f([1,r],r+1) - f([1,l),r+1).
                   [f05c5b]
                                                                                                                                                                                   Part1: Answerall f([1,r],r+1) first.
                                                                                                                                                                                   Part2: Store curR \to R for curL (reduce the space to O(N)), and then
 Mo's Algorithm With modification
                                                                                                                                                                                    answer them by the second offline algorithm.
 Block: N^{2/3}, Complexity: N^{5/3}
                                                                                                                                                                                    Note: You must do the above symmetrically for the left boundaries.
                                                                                                                                                                    9.5 Hilbert Curve [1274a3]
 struct Query {
  int L, R, LBid, RBid, T;
                                                                                                                                                                     ll hilbert(int n, int x, int y) {
       Query(int 1, int r, int t):
                                                                                                                                                                           ll res = 0;
       L(1), R(r), LBid(1 / blk), RBid(r / blk), T(t) {} bool operator<(const Query &q) const {
                                                                                                                                                                           for (int s = n / 2; s; s >>= 1) {
                                                                                                                                                                               int rx = (x \& s) > 0;
                                                                                                                                                                               int ry = (y \& s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
             \begin{array}{ll} \textbf{if} & (LBid \ != \ q.LBid) \ \textbf{return} \ LBid < q.LBid; \end{array}
             if (RBid != q.RBid) return RBid < q.RBid;</pre>
            return T < b.T;
                                                                                                                                                                                if (ry = 0) {
                                                                                                                                                                                     if (rx = 1) x = s - 1 - x, y = s - 1 - y;
                                                                                                                                                                                    swap(x, y);
 void solve(vector<Query> query) {
                                                                                                                                                                               }
                                                                                                                                                                         }
       sort(ALL(query));
       int L=0, R=0, T=-1;
                                                                                                                                                                          return res;
       for (auto q : query) { while (T < q.T) addTime(L, R, ++T); // TODO
                                                                                                                                                                     // n = 2^k 
            while (T > q.T) and Time(L, R, ++1), //
while (T > q.T) subTime(L, R, T-\cdot); //
while (R < q.R) add(arr[++R]); // TODO
while (L > q.L) add(arr[--L]); // TODO
while (R > q.R) sub(arr[R-\cdot]); // TODO
while (L < q.L) sub(arr[L++]); // TODO
                                                                                                                                                                    9.6 DynamicConvexTrick* [673ffd]
                                                                                                                                                                        / only works for integer coordinates!! maintain max
                                                                                                                                                                    struct Line {
                                                                                                                                                                          mutable ll a, b, p;
                                                                                                                                                                          bool operator
              // answer query
                                                                                                                                                                                       <(const Line &rhs) const { return a < rhs.a; }
                                                                                                                                                                          bool operator <(ll x) const { return p < x; }
}
                                                                                                                                                                    struct DynamicHull : multiset<Line, less<>> {
 9.3 Mo's Algorithm On Tree [8331c2]
                                                                                                                                                                          static const ll kInf = 1e18;
                                                                                                                                                                          ll Div(ll a,
                                                                                                                                                                                        lì b) { return a / b - ((a \hat{b}) < 0 \& a \% b); }
 Mo's Algorithm On Tree
 Preprocess:
                                                                                                                                                                          bool isect(iterator x, iterator y) {
                                                                                                                                                                                if (y == end()) { x->p = kInf; return 0; }
  1) LCA
 2) dfs with in[u] = dft++, out[u] = dft++
```

if (x

->a == y->a) x->p = x->b > y->b ? kInf : -kInf;

else x->p = Div(y->b - x->b, x->a - y->a);

= x << 13;

x = x >> 7;

x = x << 17;

return x;

```
return x->p>=y->p;
                                                                        ull dfs(int u, int f) {
   void addline(ll a, ll b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
                                                                           ull sum = seed;
                                                                           for (int i : G[u])
      while (isect(y, z)) z = erase(z);
                                                                             if (i != f)
      if (x != begin
                                                                               sum += shift(dfs(i, u));
           () && isect(--x, y)) isect(x, y = erase(y));
                                                                          return sum;
      while ((y = x) != begin
          () && (--x)-p >= y-p is ect (x, erase(y));
                                                                        9.11 Binary Search On Fraction [765c5a]
   11 query(11 x) {
    auto l = *lower_bound(x);
                                                                        struct Q {
                                                                          ll p, q;
      return l.a * x + l.b;
                                                                          Q go(Q b, 11 d) \{ return \{ p + b.p*d, q + b.q*d \}; \}
};
                                                                        bool pred(Q);
        All LCS* [78a378]
9.7
                                                                           returns smallest p/q in [lo, hi] such that
                                                                           pred(p/q) is true, and 0 \le p,q \le N
\begin{tabular}{ll} \bf void & all\_lcs(string s, string t) & {\it // 0-base} \end{tabular}
                                                                        Q frac_bs(ll N)
   vector < int > h(SZ(t));
                                                                          Q lo\{0, 1\}, hi\{1, 0\};
                                                                          if (pred(lo)) return lo;
   iota(ALL(h), 0);
   for (int a = 0; a < SZ(s); ++a) {
                                                                          assert (pred(hi));
     int v = -1;
                                                                          bool dir = 1, L = 1, H = 1;
     for (int c = 0; c < SZ(t); ++c)
if (s[a] == t[c] || h[c] < v)
                                                                          for (; L | | H; dir = !dir) {
                                                                             ll len = 0, step = 1;
                                                                             for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
          swap(h[c], v);
      \begin{array}{l} \text{if } (Q \text{ mid} = \text{hi.go}(\text{lo}\,, \text{ len} + \text{step})\,; \\ \text{mid.p} > N \mid\mid \text{mid.q} > N \mid\mid \text{dir} \, \widehat{} \text{ pred}(\text{mid})) \end{array}
      // h[i] might become -1 !!
                                                                               else len += step;
}
                                                                             swap(lo, hi = hi.go(lo, len));
                                                                             (dir ? L : H) = !!len;
        AdaptiveSimpson* [4074b3]
                                                                          return dir ? hi : lo;
template < typename Func, typename d = double >
 struct Simpson {
   using pdd = pair < d, d>;
                                                                        9.12 Bitset LCS [330ab1]
   Func f;
   pdd mix(pdd 1, pdd r, optional<d> fm = {}) {
    d h = (r.X - 1.X) / 2, v = fm.value_or(f(1.X + h));
    return {v, h / 3 * (1.Y + 4 * v + r.Y)};
                                                                        cin >> n >> m;
                                                                        for (int i = 1, x; i \le n; ++i)
                                                                          cin \gg x, p[x].set(i);
                                                                        \label{eq:formula} \mbox{for (int $i = 1$, $x$; $i <= m$; $i++$) } \{
  d eval(pdd l, pdd r, d fm, d eps) {
pdd m((l.X + r.X) / 2, fm);
                                                                          cin >> x, (g = f) |= p[x];
                                                                          f.shiftLeftByOne(), f.set(0);
     d\ s = mix(1,\ r,\ fm).second;
                                                                          ((f = g - f) = g) \& = g;
     auto [flm , sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
                                                                        cout << f.count() << '\n';
     d \ delta = sl + sr - s;
     if (abs(delta
) <= 15 * eps) return sl + sr + delta / 15;
                                                                        9.13 N Queens Problem [d1fccc]
                                                                        void solve
      return eval(1, m, flm, eps / 2) +
                                                                             (vector<int> &ret, int n) { // no sol when n=2,3
        eval(m, r, fmr, eps / 2);
                                                                           if (n % 6 == 2) {
  for (int i = 2; i <= n; i += 2) ret.pb(i);</pre>
   d eval(d l, d r, d eps) {
                                                                             ret.pb(3); ret.pb(1);
     return eval
                                                                             for (int i = 7; i \le n; i += 2) ret.pb(i);
           (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
                                                                             ret.pb(5);
  d eval2(d l, d r, d eps, int k = 997) {

d h = (r - l) / k, s = 0;

for (int i = 0; i < k; ++i, l += h)

s += eval(l, l + h, eps / k);
                                                                            else if (n \% 6 = 3) {
                                                                             for (int i = 4; i \le n; i += 2) ret.pb(i);
                                                                             ret.pb(2);
                                                                             for (int i = 5; i \le n; i += 2) ret.pb(i);
                                                                             ret.pb(1); ret.pb(3);
     return s;
                                                                            else {
   }
                                                                             for (int i = 2; i \le n; i += 2) ret.pb(i);
 };
                                                                             for (int i = 1; i \le n; i += 2) ret.pb(i);
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
9.9 Simulated Annealing [de78c6]
                                                                               Python
                                                                        10
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
                                                                        10.1
                                                                                \operatorname{Misc}
for (int it = 1; it \leq 10000000; ++it) {
                                                                        from decimal import *
                                                                        setcontext (Context (prec
           answer, nw: current value, rnd(): mt19937 rnd()
      if (exp(-(nw - ans
                                                                             =MAX_PREC, Emax=MAX_EMAX, rounding=ROUND_FLOOR))
                                                                        print(Decimal(input()) * Decimal(input()))
          ) / factor) >= (double)(rnd() % base) / base)
                                                                        from fractions import Fraction
          ans = nw;
      factor *= 0.99995;
                                                                        Fraction
                                                                             ( '3.14159 ').limit_denominator(10).numerator # 22
}
9.10 Tree Hash* [34aae5]
                                                                        \operatorname{map}(\operatorname{int},\operatorname{input}()\operatorname{.}\operatorname{split}())
                                                                        # N*M
ull seed;
                                                                        arr2d = [list]
 ull shift(ull x) {
```

(map(int,input().split()))] for i in range(N)]

" ".join(str(i) for i in a)

#a^b%M

pow(a,b,M)