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```
"This file should be placed at ~/.vimrc" se nu ai hls et ru ic is sc cul se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a syntax on hi cursorline cterm=none ctermbg=89 set bg=dark inoremap {<CR> {<CR>}<Esc>ko<tab> "Select region and then type :Hash to hash your selection." "Useful for verifying that there aren't mistypes." ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \| md5sum \| cut -c-6
```

1.2 readchar [0754b0] inline char readchar() {

```
static const size_t bufsize = 65536;
  static char buf[bufsize];
static char *p = buf, *end = buf;
  if (p == end) end = buf +
       fread_unlocked(buf, 1, bufsize, stdin), p = buf;
  return *p++;
}
1.3 Black Magic [d566f1]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace ___gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
      tree<int, null_type, std::less<int>, rb_tree_tag
      tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
      (0) = 22; assert(*s.find_by_order(1) = 71);
  assert(s.order_of_key
      (22) = 0; assert (s.order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by_order
      (0) = 71; assert (s.order_of_key(71) = 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope < char > r[2];
  r[1] = r[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0;
```

1.4 Pragma Optimization [7b330a]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse, sse2, sse3, sse3, sse4")
#pragma GCC target("popent, abm, nmmx, avx, arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

1.5 Bitset [282252]

```
#include<bits/stdc++.h>
using namespace std;

int main () {
    bitset<4> bit;
    bit.all(); // all bit is true, ret tru;
    bit.any(); // any bit is true, ret true
    bit.none(); // all bit is false, ret true
    bit.count();
    bit.to_string('0', '1');//with parmeter
    bit.reset(); // set all to true
    bit.set(); // set all to false
    std::bitset<8> b3{0}, b4{42};
    std::hash<std::bitset<8> hash_fn8(b3); hash_fn8(b4);
}
```

2 Graph

$2.1\quad BCC\ Vertex*\ [ed8308]$

```
struct BCC { // 0-base
  int n, dft, nbcc;
  vector<int> low, dfn, bln, stk, is_ap, cir;
  vector<vector<int>>> G, bcc, nG;
  void make_bcc(int u) {
    bcc.emplace_back(1, u);
    for (; stk.back() != u; stk.pop_back())
       bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
    stk.pop_back(), bln[u] = nbcc++;
```

for (; stk.back() != u; stk.pop_back())

```
bln [stk
   void dfs(int u, int f) {
                                                                                                       .back()] = nscc, instack[stk.back()] = 0;
                                                                                             instack[u] = 0, bln[u] = nscc++, stk.pop_back();
     int child = 0;
      low[u] = dfn[u] = ++dft, stk.pb(u);
      for (int v : G[u])
        if (!dfn[v]) {
                                                                                       \begin{array}{l} dfs\left(v,\;u\right),\; +\!\!\!\!+\!\!\!\!+\!\!\!\!\!+ child\;;\\ low\left[u\right]\; =\; min\left(low\left[u\right]\;,\; low\left[v\right]\right)\;; \end{array}
                                                                                       void add_edge(int u, int v) {
            if (dfn[u] \le low[v]) {
 is\_ap[u] = 1, bln[u] = nbcc;
                                                                                          G[u].pb(v);
                                                                                       make\_bcc(v), bcc.back().pb(u);
                                                                                             if (!dfn[i]) dfs(i);
        else if (dfn[v] < dfn[u] & v != f)
     low[u] = min(low[u], dfn[v]);
if (f = -1 \&\& child < 2) is ap[u] = 0;
                                                                                    }; // scc_id(i): bln[i]
     if (f = -1 & child = 0) make_bcc(u);
                                                                                     2.4 2SAT* [e839e5]
   \begin{array}{lll} \dot{B}CC(int \ \underline{\ \ } n): \ n(\underline{\ \ \ } n) \ , \ dft \, () \ , \\ nbcc \, () \ , \ low \, (n) \ , \ dfn \, (n) \ , \ bln \, (n) \ , \ is\_ap \, (n) \ , \ G(n) \ \ \{\} \\ void \ add\_edge \, (int \ u, \ int \ v) \ \ \{ \end{array} 
                                                                                    struct SAT { // 0-base
                                                                                       int n:
                                                                                       vector<bool> istrue;
     G[u].pb(v), G[v].pb(u);
                                                                                       SCC scc;
                                                                                       SAT(int \underline{n}): n(\underline{n}), istrue(n+n), scc(n+n)  {}
  void solve() {
  for (int i = 0; i < n; ++i)</pre>
                                                                                       int rv(int a) {
                                                                                          return a > = n ? a - n : a + n;
         if (!dfn[i]) dfs(i, -1);
                                                                                        void add_clause(int a, int b) {
   void block_cut_tree() {
                                                                                          scc.add\_edge(rv(a), b), scc.add\_edge(rv(b), a);
      cir.resize(nbcc);
      for (int i = 0; i < n; ++i)
        if (is_ap[i])
bln[i] = nbcc++;
                                                                                       bool solve()
                                                                                          scc.solve();
                                                                                           for (int i = 0; i < n; ++i) {
     cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
                                                                                             if(scc.bln[i] = scc.bln[i+n]) return false;
                                                                                             istrue[i] = scc.bln[i] < scc.bln[i + n];
         for (int j : bcc[i])
                                                                                             istrue[i + n] = !istrue[i];
            if (is_ap[j])
  nG[i].pb(bln[j]), nG[bln[j]].pb(i);
} // up to 2 * n - 2 nodes!! bln[i] for id
                                                                                           return true;
                                                                                       }
                                                                                    };
2.2 Bridge* [Occada]
                                                                                            MinimumMeanCycle* [4be648]
struct ECC { // 0-base
   int n, dft, ecnt, necc;
vector<int> low, dfn, bln, is_bridge, stk;
                                                                                     ll road[N][N]; // input here
                                                                                     struct MinimumMeanCycle {
   {\tt vector}{<\hspace{-1.5pt}{\rm vector}}{<\hspace{-1.5pt}{\rm pii}}{>\hspace{-1.5pt}{>\hspace{-1.5pt}{>}}}\;G;
                                                                                       11 \, dp[N + 5][N], n;
                                                                                        pll solve() {
   void dfs(int u, int f) {
                                                                                          11 a = -1, b = -1, L = n + 1;
for (int i = 2; i \le L; ++i)
      dfn\left[u\right] \,=\, low\left[u\right] \,=\, +\!\!+\!\! dft \;,\;\; stk \,.\, pb(u) \;;
     for (auto [v, e] : G[u])
if (!dfn[v])
                                                                                             for (int k = 0; k < n; ++k)
         \begin{array}{l} dfs\left(v,\,e\right),\;low\left[u\right] \,=\, min(low\left[u\right],\;low\left[v\right])\,;\\ else\;\;if\;\left(e\;!=\;f\right) \end{array}
                                                                                                for (int j = 0; j < n; ++j)
                                                                                                   dp[i][j] =
                                                                                           \min_{\substack{m \text{ in } (dp[i-1][k] + road[k][j], \ dp[i][j]); \\ \text{for } (int \ i=0; \ i< n; ++i) \ \{ } 
     \begin{array}{c} low\left[u\right] \stackrel{.}{=} min(low\left[u\right],\ dfn\left[v\right])\,;\\ if\ (low\left[u\right] \stackrel{.}{==} dfn\left[u\right])\ \{ \end{array}
                                                                                             if (dp[L][i] >= INF) continue;
         if (f != -1) is_bridge[f] = 1;
                                                                                             for (; stk.back() != u; stk.pop_back())
bln[stk.back()] = necc;
        bln[u] = necc++, stk.pop\_back();
                                                                                                   ta \, = \, \dot{d}p \, [L] \, [\, \dot{i}\, ] \  \, - \, dp \, [\, \dot{j}\, ] \, [\, \dot{i}\, ] \, , \  \, tb \, = \, L \, \, - \, \, \dot{j} \, ;
                                                                                             if (ta == 0) continue;
  ÉCC(int _n): n(_n), dft()
         , \ \underline{ecnt}() \ , \ \underline{necc}() \ , \ low(\underline{n}) \ , \ dfn(\underline{n}) \ , \ bln(\underline{n}) \ , \ G(\underline{n}) \ \{\}
                                                                                             if (a = -1) | a * tb > ta * b) a = ta, b = tb;
   void add_edge(int u, int v)
                                                                                          if (a != -1) \{

11 g = gcd(a, b);
     G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
                                                                                             return pll(a / g, b / g);
   void solve() {
     is_bridge.resize(ecnt);
for (int i = 0; i < n; ++i)</pre>
                                                                                          return pll(-1LL, -1LL);
         if (!dfn[i]) dfs(i, -1);
                                                                                       void init(int _n) {
}; // ecc_id(i): bln[i]
                                                                                           for (int i = 0; i < n; ++i)
2.3 SCC* [22afe1]
                                                                                              \mbox{for (int } j \, = \, 0 \, ; \ j \, < \, n \, ; \ + + j \, ) \ dp [\, i \, + \, 2 \, ][\, j \, ] \, = \, INF \, ; 
struct SCC { // 0-base
                                                                                    };
  int n, dft, nscc;
   \label{eq:vector} vector <\!\!int\!\!> low\,, \ dfn\,, \ bln\,, \ instack\,, \ stk\,;
                                                                                    2.6 Virtual Tree* [80f7cb]
   vector < vector < int >>> G;
   void dfs(int u)
                                                                                     vector < int > vG[N];
     low[u] = dfn[u] = ++dft;
                                                                                     int top, st[N];
     instack[u] = 1, stk.pb(u);
     for (int v : G[u])
                                                                                     void insert(int u) {
         if^{(dfn[v])}
                                                                                       if (top = -1) return st[++top] = u, void();
         d\hat{f}s(v), low[u] = min(low[u], low[v]);
else if (instack[v] && dfn[v] < dfn[u])
                                                                                       int p = LCA(st[top], u);
                                                                                        if (p = st[top]) return st[++top] = u, void();
                                                                                       while (top >= 1 \&\& dep[st[top - 1]] >= dep[p])
           low[u] = min(low[u], dfn[v]);
      if (low[u] = dfn[u]) {
                                                                                        vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
```

dst[i][i] = vcst[i] = 0;

```
void chmin(int &x, int val) {
     vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
                                                                                   x = \min(x, val);
                                                                                 void add_edge(int ui, int vi, int wi) {
void reset(int u) {
                                                                                   chmin (dst [ui][vi], wi);
  for (int i : vG[u]) reset(i);
  vG[u].clear();
                                                                                void shortest_path() {
                                                                                   for (int k = 0; k < n; ++k)
                                                                                     for (int i = 0; i < n; ++i)
                                                                                        for (int j = 0; j < n; ++j)
void solve(vector<int> &v) {
                                                                                           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  top = -1
  sort (ALL(v),
     [\&](int a, int b) \{ return dfn[a] < dfn[b]; \});
                                                                                int solve(const vector<int>& ter) {
  for (int i : v) insert(i);
                                                                                   shortest_path();
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
                                                                                   int t = SZ(ter), full = (1 << t) - 1;
  // do something
                                                                                   for (int i = 0; i \leftarrow full; ++i)
                                                                                     fill_n (dp[i], n, INF);
  reset(v[0]);
                                                                                   \begin{array}{l} \operatorname{copy\_n}(\operatorname{vcst},\ \operatorname{n},\ \operatorname{dp}[\ 0\ ])\ ;\\ \operatorname{for}\ (\operatorname{int}\ \operatorname{msk}=\ 1;\ \operatorname{msk}<=\ \operatorname{full}\ ;\ +\!\!+\!\!\operatorname{msk})\ \{ \end{array}
       Maximum Clique Dyn* [4a6b3d]
                                                                                     if (!(msk & (msk - 1))) {
                                                                                        \begin{array}{l} \text{int who} = \underline{\hspace{0.5cm}} \lg(msk); \\ \text{for (int } i = 0; i < n; ++i) \end{array}
struct MaxClique { // fast when N \le 100 bitset A > G[N], cs[N]; int ans, sol[N], q, cur[N], d[N], n; void init (int _n) {
                                                                                           dp [msk
                                                                                                [i] = vcst[ter[who]] + dst[ter[who]][i];
     n = n;
                                                                                      for (int i = 0; i < n; ++i)
      \  \, \text{for} \  \, (\, \text{int} \  \, i \, = \, 0\,; \  \, i \, < \, n\,; \, +\!\!\!\!+\!\!\! i\,) \, \, G[\, i\, ]\,.\, reset\,(\,)\,; \\
                                                                                        for (int sub = (
                                                                                             msk - 1) & msk; sub; sub = (sub - 1) & msk)
  void add_edge(int u, int v) {
                                                                                           \begin{array}{c} chmin(dp[msk][i],\\ dp[sub][i] + dp[msk \ \widehat{\ } sub][i] - vcst[i]); \end{array}
    G[u][v] = G[v][u] = 1;
                                                                                     for (int i = 0; i < n; ++i) { tdst[i] = INF;
  \begin{tabular}{ll} void $$ pre\_dfs(vector<int> \&r, int 1, bitset< N> mask) $$ \{$ \end{tabular}
     if (1 < 4) {
                                                                                         for (int j = 0; j < n; ++j)
        for (int i : r) d[i] = (G[i] & mask).count();
                                                                                           chmin(tdst[i], dp[msk][j] + dst[j][i]);
        sort (ALL(r)
             , [\&](int x, int y) \{ return d[x] > d[y]; \});
                                                                                     copy_n(tdst, n, dp[msk]);
     vector < int > c(SZ(r));
                                                                                   return *min_element(dp[full], dp[full] + n);
     int 1 \text{ft} = \max(\text{ans} - q + 1, 1), \text{ rgt} = 1, \text{ tp} = 0;
     cs[1].reset(), cs[2].reset();
                                                                             ; // O(V 3^T + V^2 2^T)
     for (int p : r) {
                                                                              2.9 Dominator Tree* [e95beb]
        int k = 1;
        while ((cs[k] \& G[p]).any()) ++k;
        if (k > rgt) cs[++rgt + 1].reset();
cs[k][p] = 1;
if (k < lft) r[tp++] = p;</pre>
                                                                              struct dominator_tree
                                                                                                              // 1-base
                                                                                vector<int> G[N], rG[N]; int n, pa[N], idn[N], id [N], Time; int semi[N], idom[N], best[N]; vector<int> tree[N]; // dominator_tree
     for (int k = lft; k \ll rgt; ++k)
                                                                                 void init(int _n) {
        for (int p = cs[k]._Find_first
                                                                                  n = _n;
for (int i = 1; i <= n; ++i)
rG[i]. clear(
          (); p < N; p = cs[k]._Find_next(p))
r[tp] = p, c[tp] = k, ++tp;
                                                                                     G[i].clear(), rG[i].clear();
     dfs(r, c, l + 1, mask);
                                                                                 void add_edge(int u, int v) {
  void dfs (vector<
                                                                                  G[u].pb(v), rG[v].pb(u);
        int>&r, vector<int>&c, int 1, bitset<N> mask) {
     while (!r.empty()) {
                                                                                void dfs(int u)
        int p = r.back();
                                                                                   id[dfn[u] = ++Time] = u;
        r.pop\_back(), mask[p] = 0;
                                                                                   for (auto v : G[u])
        if (q + c.back() <= ans) return;</pre>
                                                                                     if(!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
        \operatorname{cur}\left[q++\right] = p;
        vector<int> nr;
                                                                                int find(int y, int x) {
        for (int i : r) if (G[p][i]) nr.pb(i);
                                                                                   if (y \le x) return y;
        if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
                                                                                   int tmp = find(pa[y], x);
        else if (q > ans) ans = q, copy_n(cur, q, sol); c.pop_back(), --q;
                                                                                   if (semi[best[y]] > semi[best[pa[y]]])
                                                                                     best[y] = best[pa[y]];
     }
                                                                                   return pa[y] = tmp;
  int solve() {
                                                                                void tarjan(int root) {
     vector < int > r(n);
                                                                                  Time = 0;
     ans = q = 0, iota(ALL(r), 0);
                                                                                   for (int i = 1; i \le n; ++i) { dfn[i] = idom[i] = 0;
     pre\_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
                                                                                      tree[i].clear();
                                                                                      best[i] = semi[i] = i;
};
                                                                                   dfs(root);
for (int i = Time; i > 1; --i) {
        Minimum Steiner Tree* [cbf811]
struct SteinerTree {
                                                                                     int u = id[i];
for (auto v : rG[u])
                            // 0-base
  int n, dst[N][N], dp[1 \ll T][N], tdst[N]; int vcst[N]; // the cost of vertexs
                                                                                        \inf (v = dfn[v]) \{
                                                                                           find(v, i);

semi[i] = min(semi[i], semi[best[v]]);
  void init(int _n) {
     for (int i = 0; i < n; ++i) {
        fill_n(dst[i], n, INF);
                                                                                      tree [semi[i]].pb(i);
```

for (auto v: tree[pa[i]]) {

find(v, pa[i]);idom[v] =

```
semi\left[\,best\left[\,v\,\right]\,\right] \;=\!\!\!-\; pa\left[\,i\,\right] \;\;?\;\; pa\left[\,i\,\right] \;\; : \;\; best\left[\,v\,\right];
          tree [pa[i]]. clear();
      for (int i = 2; i \leftarrow Time; ++i) {
          if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
          tree [id [idom[i]]].pb(id[i]);
};
```

Minimum Clique Cover* [5951ca] 2.10

```
\begin{array}{lll} & \textbf{struct} & Clique\_Cover \; \{ \; \; // \; 0\text{-base} \,, \; O(n2 \hat{}^{}n) \\ & \textbf{int} \; \; co[\; 1 << N] \,, \; n, \; E[N] \,; \\ & \textbf{int} \; \; dp[\; 1 << N] \,; \end{array}
    void init (int _n) {
    n = _n, fill_n (dp, 1 << n, 0);

        fill_n(E, n, 0), fill_n(co, 1 << n, 0);
    int solve() {
  for (int i = 0; i < n; ++i)
            co[1 << i] = E[i] | (1 << i);
        co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
        for (int i = 1; i < (1 << n); ++i) {
            or (int i -
int t = i & -i;
dn[i ^ t];
            \begin{array}{l} dp[\,i\,] \,=\, -dp[\,i\,\,\widehat{\ }\,\, t\,\,]\,;\\ co\,[\,i\,\,] \,=\, co\,[\,i\,\,\widehat{\ }\,\, t\,\,]\,\,\&\,\, co\,[\,t\,\,]\,; \end{array}
        for (int i = 0; i < (1 << n); ++i)

co[i] = (co[i] & i) == i;
         fwt(co, 1 << n, 1);
        for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)
    sum += (dp[i] *= co[i]);</pre>
             if (sum) return ans;
        return n;
    }
};
```

NumberofMaximalClique* [c163d7]

```
struct BronKerbosch { // 1-base
   \begin{array}{l} int \ n, \ a\left[N\right], \ g\left[N\right]\left[\stackrel{.}{N}\right]^{'}, \\ int \ S, \ all \left[N\right]\left[\stackrel{.}{N}\right], \ some \left[N\right]\left[N\right], \ none \left[N\right]\left[N\right]; \end{array}
    void init(int _n) {
       n = \underline{n};
       for (int i = 1; i \le n; ++i)
           for (int j = 1; j \le n; ++j) g[i][j] = 0;
    void add_edge(int u, int v) {
       g[u][v] = g[v][u] = 1;
    void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
       if (sn = 0 \&\& nn = 0) + S;
       int u = some[d][0];
       for (int i = 0; i < sn; ++i) {
           int v = some[d][i];
          int v = some[d][1];
if (g[u][v]) continue;
int tsn = 0, tnn = 0;
copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
for (int j = 0; j < sn; ++j)
    if (g[v][some[d][j]])</pre>
           some [d + 1][tsn++] = some[d][j];
for (int j = 0; j < nn; ++j)
               if (g[v][none[d][j]])
                  none[d + 1][tnn++] = none[d][j];
           \begin{array}{l} dfs\,(d+1,\;an+1,\;tsn\,,\;tnn)\,;\\ some\,[d]\,[\,i\,]\,=\,0\,,\;none\,[d\,]\,[\,nn++]\,=\,v\,; \end{array}
    int solve() {
       iota(some[0], some[0] + n, 1);
       S = 0, dfs(0, 0, n, 0);
       return S;
};
```

3 Data Structure

```
3.1 Discrete Trick
vector<int> val;
// build
sort (ALL
    (val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// \max idx <= :
upper_bound(ALL(val), x) - val.begin();
// \max idx < x
lower_bound(ALL(val), x) - val.begin();
3.2 BIT kth* [7d1b5f]
int bit [N + 1]; // N = 2 ^ k
int query_kth(int k) {
    int res = 0;
    for (int i = N >> 1; i >= 1; i >>= 1)
        if (bit [res + i] < k)
            k -= bit [res += i];
    return res + 1;
3.3 Interval Container* [78516e]
/* Add and
     remove intervals from a set of disjoint intervals.
 * Will merge the added interval with
     any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
\mathtt{set} \! < \! \mathtt{pii} > ::
    iterator addInterval(set<pii> is, int L, int R) {
  if (L == R) return is.end();
```

$R = \max(R, \text{ it ->Y});$

```
auto it = is.lower_bound(\{L, R\}), before = it; while (it != is.end() && it->X <= R) {
  before = it = is.erase(it);
if (it != is.begin() && (--it)->Y>= L) {
  L = \min(L, it ->X);
  R = \max(R, it ->Y);
  is.erase(it);
```

```
void removeInterval(set<pii> is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it ->Y;
  if (it->X == L) is.erase(it);
  else (int\&)it ->Y = L;
  if (R != r2) is .emplace(R, r2);
```

return is.insert(before, pii(L, R));

3.4 Leftist Tree [bbd228]

```
ll v, data, sz, sum; node *1, *r;
  node(ll k)
       v(0), data(k), sz(1), l(0), r(0), sum(k) {}
ll sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
11 node *merge(node *a, node *b) {
  if (!a | | '!b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
  delete tmp;
```

3.5 Heavy light Decomposition* [babe8a]

```
struct Heavy_light_Decomposition {
  int n, ulink [N], deep [N], mxson [N], w[N], pa [N];
  int t, pl[N], data[N], val[N]; // val: vertex data
```

```
National Taiwan University 8BQube
   \begin{array}{l} {\rm vector}{<} {\rm int}{>} \; G[N] \, ; \\ {\rm void \;\; init} \; ({\rm int} \;\; \underline{\ \ \ \ \ } ) \;\; \{ \end{array}
     n = n;
for (int i = 1; i <= n; ++i)
        G[i].clear(), mxson[i] = 0;
   void add_edge(int a, int b) {
     G[a].pb(b), G[b].pb(a);
   void dfs(int u, int f, int d) {
     w[u] = 1, pa[u] = f, deep[u] = d++; for (int &i : G[u])
         if (i != f) {
           dfs(i, u, d), w[u] += w[i];
            if (w[mxson[u]] < w[i]) mxson[u] = i;
    \begin{array}{lll} void & cut(int \ u, \ int \ link) \ \{ \\ & data[pl[u] = +\!\!+\!\!t] = val[u] \ , \ ulink[u] = link \ ; \end{array} 
      if (!mxson[u]) return;
                                                                                      }
      cut(mxson[u], link);
for (int i : G[u])
                                                                                   };
         if (i != pa[u] && i != mxson[u])
           cut(i, i);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
      \begin{array}{ll} \mbox{int } ta = \mbox{ulink} [a] \,, \ tb = \mbox{ulink} [b] \,, \ res = 0; \\ \mbox{while} \ (ta := tb) \, \{ \end{array}
         if (deep
          [ta] > deep[tb]) swap(ta, tb), swap(a, b);

/ query(pl[tb], pl[b])
         tb = ulink[b = pa[tb]];
      if (pl[a] > pl[b]) swap(a, b);
      // query(pl[a], pl[b])
        Centroid Decomposition* [4eccaf]
void init(int _n) {
     n = \underline{n}, layer[0] =
                                 -1;
      fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
      for (int i = 1; i \le n; ++i) G[i]. clear();
```

```
void add_edge(int a, int b, int w) {
   G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
void get_cent(
  int u, int f, int &mx, int &c, int num) {
   int mxsz = 0;
   sz[u] = 1;
   for (pll e : G[u])
      if (!done[e.X] && e.X != f) {
  get_cent(e.X, u, mx, c, num);
         sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
   if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
void dfs(int u, int f, ll d, int org) {
  // if required, add self info or climbing info
  dis[layer[org]][u] = d;
  for (pll e : G[u])
    if (!done[e.X] && e.X != f)
        dfs(e.X, u, d + e.Y, org);
         dfs(e.X, u, d + e.Y, org);
int cut(int u, int f, int num) {
   int mx = 1e9, c = 0, lc;
get_cent(u, f, mx, c, num);
   done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1; for (pll e : G[c])
      if (!done[e.X]) {
         if (sz[e.X] > sz[c])
  lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
         upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
   return done[c] = 0, c;
```

```
void build() { cut(1, 0, n); }
void modify(int u) {
   for (int a = u, ly = layer[a]; a;
        a = pa[a], --ly) {
        info[a].X += dis[ly][u], ++info[a].Y;
        if (pa[a])
            upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
}

ll query(int u) {
    ll rt = 0;
   for (int a = u, ly = layer[a]; a;
        a = pa[a], --ly) {
        rt += info[a].X + info[a].Y * dis[ly][u];
        if (pa[a])
        rt -=
            upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    }
   return rt;
}
```

3.7 LiChaoST* [4a61ec]

```
struct L {
  ll m, k, id;
  \begin{array}{l} L() \ : \ id(-1) \ \{\} \\ L(11 \ a, \ 11 \ b, \ 11 \ c) \ : \ m(a) \, , \ k(b) \, , \ id(c) \ \{\} \\ 11 \ at(11 \ x) \ \{ \ return \ m \ ^* \ x + k; \ \} \end{array}
class LiChao { // maintain max
private:
  int n; vector<L> nodes;
  if (nodes[rt].id = -1)
        return nodes[rt] = ln, void();
     bool atLeft = nodes[rt].at(1) < ln.at(1);
     if (nodes[rt].at(m) < ln.at(m))
     atLeft = 1, swap(nodes[rt], ln);
if (r - l == 1) return;
      if \ (atLeft) \ insert(l\,,\,m,\ rt << 1\,,\,ln); \\
     else insert(m, r, rt \ll 1 | 1, ln);
  Il query(int 1, int r, int rt, ll x) {
  int m = (l + r) >> 1; ll ret = -INF;
  if (nodes[rt].id != -1) ret = nodes[rt].at(x);
  if (r - l == 1) return ret;
     public:
  LiChao(int n_) : n(n_), nodes(n * 4) {}
  void insert(\overline{L} ln) { insert(0, n, 1, ln);
  11 query(11 x) { return query(0, n, 1, x); }
```

3.8 Link cut tree* [5f036a]

```
struct Splay { // xor-sum
  static Splay nil;
  Splay *ch[2], *f;
   int val, sum, rev, size;
   Splay (int
   val = 0) : val(val), sum(val), rev(0), size(1) { f = ch[0] = ch[1] = &nil; }
   bool isr()
   { return f->ch[0] != this && f->ch[1] != this; }
   int dir()
   { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
     ch[d] = c;
      if (c != \&nil) c -> f = this;
     pull();
   void give_tag(int r) {
     if (r) swap(ch[0], ch[1]), rev = 1;
   void push() {
     if (ch[0] != &nil) ch[0]->give_tag(rev);
if (ch[1] != &nil) ch[1]->give_tag(rev);
     rev = 0;
   void pull() {
      // take care of the nil!
      size = ch[0] - size + ch[1] - size + 1;
```

function<bool(const point &, const point &)> f =
 [dep](const point &a, const point &b) {
 if (dep & 1) return a.x < b.x;</pre>

```
else return a.y < b.y;
     if (\operatorname{ch}[1] != \& \operatorname{nil}) \operatorname{ch}[1] -> f = \operatorname{this};
                                                                                   int m = (l + r) >> 1;

nth\_element(p + l, p + m, p + r, f);
                                                                                   xl[m] = xr[m] = p[m].x;
} Splay::nil;
Splay : nil;
Splay *nil = &Splay :: nil;
void rotate(Splay *x) {
                                                                                   yl[m] = yr[m] = p[m] \cdot y;

lc[m] = build(1, m, dep + 1);
void rotate (Splay
  Splay *p = x \rightarrow f;
                                                                                   if (~lc[m]) {
   int d = x - sdir();
                                                                                     xl[m] = min(xl[m], xl[lc[m]]);
                                                                                     \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{lc}[m]]);
   if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
p->setCh(x->ch[!d], d);
                                                                                     yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
  x->setCh(p, !d);
  p->pull(), x->pull();
                                                                                   rc[m] = build(m + 1, r, dep + 1);
                                                                                   if (~rc[m]) {
void splay (Splay *x) {
                                                                                     xl[m] = min(xl[m], xl[rc[m]]);
  vector < Splay * splay Vec;
for (Splay * q = x;; q = q->f) {
                                                                                     xr[m] = max(xr[m], xr[rc[m]]);
                                                                                     yl[m] = min(yl[m], yl[rc[m]]);
     splayVec.pb(q);
                                                                                     yr[m] = max(yr[m], yr[rc[m]]);
     if (q->isr()) break;
                                                                                   return m;
   reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
while (!x->isr()) {
                                                                                bool bound (const point &q, int o, long long d) {
                                                                                   \begin{array}{l} \mbox{double ds} = \mbox{sqrt}(d+1.0); \\ \mbox{if } (q.x < xl[o] - ds \mid\mid q.x > xr[o] + ds \mid\mid \\ \mbox{q.y} < yl[o] - ds \mid\mid q.y > yr[o] + ds) \end{array} 
     if (x->f->isr()) rotate(x);
     else if (x->dir() = x->f->dir())
        rotate(x->f), rotate(x);
                                                                                     return false;
     else rotate(x), rotate(x);
                                                                                   return true;
                                                                                long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
Splay* access(Splay *x) {
  Splay *q = nil;
for (; x != nil; x = x->f)
     \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
                                                                                   const point &q, long long &d, int o, int dep = 0) {
   return q;
                                                                                   if (!bound(q, o, d)) return;
long long cd = dist(p[o], q);
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
                                                                                   if (cd != 0) d = min(d, cd);
  root_path(x), x->give_tag(1);
                                                                                   if ((\text{dep \& 1}) \&\& q.x < p[o].x | |
                                                                                     !(dep & 1) && q.y < p[o].y) {
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  x->push(), x->pull();
void split (Splay *x, Splay *y) {
  chroot(x), root_path(y);
                                                                                   } else {
                                                                                      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void link (Splay *x, Splay *y) {
  root_path(x), chroot(y);
x->setCh(y, 1);
                                                                                  }
                                                                                void init(const vector<point> &v) {
                                                                                   for (int i = 0; i < v.size(); ++i) p[i] = v[i];
void cut(Splay *x, Splay *y) {
                                                                                   root = build(0, v.size());
   split(x, y);
   if (y->size != 5) return;
  y->push();
                                                                                long long nearest (const point &q) {
  y->ch[0] = y->ch[0]->f = nil;
                                                                                   long long res = 1e18;
                                                                                   dfs(q, res, root);
Splay* get_root(Splay *x)
                                                                                   return res;
   for (root_path(x); x->ch[0] != nil; x = x->ch[0])
     x \rightarrow push();
                                                                                  // namespace kdt
   splay(x);
                                                                                      Flow/Matching
  return x;
                                                                                4.1 Bipartite Matching* [f07280]
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
                                                                                \begin{array}{lll} \textbf{struct} & \textbf{Bipartite\_Matching} ~\{~ / / ~ \textbf{0-base} \\ & \textbf{int} & \textbf{mp[N]} ~, ~ \textbf{mq[N]} ~, ~ \textbf{dis} [N + 1] ~, ~ \textbf{cur} [N] ~, ~ 1 ~, ~ r ~; \end{array}
Splay* lca(Splay *x, Splay *y) {
   access(x), root_path(y);
                                                                                   vector < int > G[N + 1];
                                                                                   bool dfs(int u) {
   if (y->f = nil) return y;
                                                                                      for (int &i = cur[u]; i < SZ(G[u]); ++i) {
   return y->f;
                                                                                        int e = G[u][i];
                                                                                         if (mq[e]
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
                                                                                                (\operatorname{dis}[\operatorname{mq}[e]] = \operatorname{dis}[u] + 1 \& \operatorname{dfs}(\operatorname{mq}[e]))
                                                                                           return mp[mq[e] = u] = e, 1;
int query (Splay *x, Splay *y) {
                                                                                     return dis[u] = -1, 0;
  split(x, y);
  return y->sum;
                                                                                   bool bfs() {
                                                                                     queue<int> q;
                                                                                      fill_n(dis, l + 1, -1);
3.9 KDTree [74016d]
                                                                                      for (int i = 0; i < l; ++i)
                                                                                        if (!~mp[i])
namespace kdt {
int root, lc [maxn], rc [maxn], xl [maxn], xr [maxn],
                                                                                          q.push(i), dis[i] = 0;
                                                                                     while (!q.empty()) {
  yl[maxn], yr[maxn];
point p[maxn];
                                                                                        int u = q.front();
int build(int l, int r, int dep = 0) {
                                                                                        q.pop();
   if (l = r) return -1;
```

if (!~dis[mq[e]])

q.push(mq[e]), dis[mq[e]] = dis[u] + 1;

```
return dis[1] != -1;
}
int matching() {
  int res = 0;
  fill_n(mp, 1, -1), fill_n(mq, r, 1);
  while (bfs()) {
    fill_n(cur, 1, 0);
    for (int i = 0; i < 1; ++i)
        res += (!~mp[i] && dfs(i));
    }
  return res; // (i, mp[i] != -1)
}
void add_edge(int s, int t) { G[s].pb(t); }
void init(int_l, int_r) {
    l = _l, r = _r;
    for (int i = 0; i <= 1; ++i)
        G[i].clear();
}
}</pre>
```

4.2 Kuhn Munkres* [edf909]

```
{\tt struct} KM { // 0-base, maximum matching
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)
      fill_n (w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
  if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)
         for (int x = 0, y = qu[ql++]; x < n; ++x)
           if (!vl[x] && slk
                [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d; else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x])' d = slk[x];
       for (int x = 0; x < n; ++x) {
         \begin{array}{l} \text{if } (vl[x]) \ hl[x] += d; \\ \text{else } slk[x] -= d; \end{array}
         if (\operatorname{vr}[x]) \operatorname{hr}[x] = d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
  ll solve() {
    fill_n(fl
         , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
    for (int i = 0; i < n; ++i)
    hl[i] = *max\_element(w[i], w[i] + n);
for (int i = 0; i < n; ++i) bfs(i);
    ll res = 0;
    return res;
};
```

4.3 MincostMaxflow* [47d9d2]

```
struct MinCostMaxFlow { // 0-base
    struct Edge {
        1l from, to, cap, flow, cost, rev;
    } *past[N];
    vector<Edge> G[N];
    int inq[N], n, s, t;
    ll dis[N], up[N], pot[N];
    bool BellmanFord() {
        fill_n(dis, n, INF), fill_n(inq, n, 0);
    }
}
```

```
queue<int> q;
auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if \ (cap > 0 \ \&\& \ dis [u] > d) \ \{ \\
       dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
  relax(s, 0, INF, 0);
   while (!q.empty())
    int u = q.front();
     q.pop(), inq[u] = 0;
for (auto &e : G[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
             (e.to, d2, min(up[u], e.cap - e.flow), &e);
    }
  }
  return dis[t] != INF;
void solve (int
  , int _t, ll &flow, ll &cost, bool neg = true) { s = \_s, t = \_t, flow = 0, cost = 0;
   if (neg) BellmanFord(), copy_n(dis, n, pot);
   for (; BellmanFord(); copy_n(dis, n, pot)) {
     for (int
     i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
     for (int i = t; past[i]; i = past[i]->from) {
  auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
  }
}
void init(int _n) {
    n = _n, fill_n(pot, n, 0);
  for (int i = 0; i < n; ++i) G[i].clear();
void add_edge(ll a, ll b, ll cap, ll cost)
  \begin{array}{l} G[a].p\overline{b}(Edge\{a,\ b,\ cap\ ,\ 0,\ cost\ ,\ SZ(G[b])\ ))\,;\\ G[b].pb(Edge\{b,\ a,\ 0,\ 0,\ -cost\ ,\ SZ(G[a])\ -\ 1\}); \end{array}
```

4.4 Maximum Simple Graph Matching* [233755]

```
struct Matching { // 0-base
   queue < int > q; int n;
   vector<int> fa, s, vis, pre, match;
   vector<vector<int>>> G;
   int Find(int u)
   \{ return u = fa[u] ? u : fa[u] = Find(fa[u]); \}
   int LCA(int x, int y) {
      static int tk = 0; tk++; x = Find(x); y = Find(y); for (;; swap(x, y)) if (x != n) {
    if (vis[x] == tk) return x;
         vis[x] = tk;
         x = Find(pre[match[x]]);
   void Blossom(int x, int y, int l) {
  for (; Find(x) != l; x = pre[y]) {
         pre[x] = y, y = match[x];
         if (s[y] = 1) q.push(y), s[y] = 0;
for (int z: \{x, y\}) if (fa[z] = z) fa[z] = 1;
   bool Bfs(int r) {
      \begin{array}{ll} iota\left(ALL(fa),\ 0\right); & fill\left(ALL(s),\ -1\right); \\ q = queue < int > (); & q.push(r); & s[r] = 0; \end{array}
      for (; !q.empty(); q.pop()) {
         for (int x = q.front(); int u : G[x])
             if(s[u] = -1) {
                if (pre[u] = x, s[u] = 1, match[u] == n) {
                   for (int a = u, b = x, last;
                        b != n; a = last, b = pre[a]
                      last =
                            match[b], match[b] = a, match[a] = b;
                  return true:
              \begin{array}{l} \text{q.push}(\mathrm{match}[u])\,; \ s[\mathrm{match}[u]] = 0\,; \\ \} \ else \ \ \text{if} \ (!\,s[u] \ \&\& \ \mathrm{Find}(u) \ != \ \mathrm{Find}(x)) \ \{ \end{array} 
               int l = LCA(u, x);
Blossom(x, u, l); Blossom(u, x, l);
```

return false;

```
National Taiwan University 8BQube
   Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis
   \begin{array}{c} (n+1)\,,\; pre\,(n+1,\;n)\,,\; match\,(n+1,\;n)\,,\; G(n)\ \{\}\\ void\ add\_edge\,(int\ u,\; int\ v) \end{array}
    \left\{ \begin{array}{l} G[u].pb(v), G[v].pb(u); \\ int \ solve() \end{array} \right\} 
      int ans = 0;
      for (int x = 0; x < n; ++x)
         if (match[x] == n) ans += Bfs(x);
      return ans;
  } // match[x] = n means not matched
       Maximum Weight Matching* [c80005]
#define REP(i, l, r) for (int i=(l); i<=(r); ++i) struct WeightGraph { // 1-based
   struct edge { int u, v, w; }; int n, nx;
   vector<int> lab; vector<vector<edge>>> g;
   vector<int> slk, match, st, pa, S, vis;
  vector<int> six, match, st, pa, b, vis, vector<vector<int> flo, flo_from; queue<int> q; WeightGraph(int n_): n(n_), nx(n * 2), lab(nx + 1), g(nx + 1, vector<edge>(nx + 1)), slk(nx + 1),
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
      match = st = pa = S = vis = slk;
     \label{eq:rep} \text{REP}(u, \ 1, \ n) \ \hat{\text{REP}}(v, \ 1, \ n) \ g[u][v] = \{u, \ v, \ 0\};
   int E(edge e)
   { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; } void update_slk(int u, int x, int &s)
   { if (!s | E(g[u][x]) < E(g[s][x]))'s = u; } void set_slk(int x) {
     slk[x] = 0;

REP(u, 1, n)
         if (g[u][x].w > 0 \& st[u] != x \& S[st[u]] == 0)
            update_slk(u, x, slk[x]);
   void q_push(int x)
      if (x \le n) q.push(x);
      else for (int y : flo[x]) q_push(y);
   void set_st(int x, int b) {
      st[x] = b;
      if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
         reverse(1 + ALL(f)), it = f.end() - pr;
      auto res = vector(f.begin(), it);
                                                                                                       }
      return f.erase(f.begin(), it), res;
   void set_match(int u, int v) {
  match[u] = g[u][v].v;
      if (u \le n) return;
      \begin{array}{l} \text{int } xr = flo\_from[u][g[u][v].u]; \\ \text{auto } \&f = flo[u], \ z = split\_flo(f, \ xr); \\ \text{REP}(i, \ 0, SZ(z) - 1) \ \text{set\_match}(z[i], \ z[i \ \cap \ 1]); \\ \end{array} 
      set_match(xr, v); f.insert(f.end(), ALL(z));
   void augment(int u, int v) {
      for (;;) {
         int xnv = st[match[u]]; set_match(u, v);
         if (!xnv) return;
         set_match(v = xnv, u = st[pa[xnv]]);
   int lca(int u, int v) {
      static int t = 0; ++\hat{t};
      for (++t; u | | v; swap(u, v)) if (u) {
  if (vis[u] == t) return u;
  vis[u] = t, u = st[match[u]];
         if (u) u = st[pa[u]];
      return 0;
   void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0, match[b] = match[o];
      vector < int > f = \{o\};
```

for (int t : {u, v}) {
 reverse(1 + ALL(f));

 $flo[b] = f; set_st(b, b);$

for (int x = t, y; x != o; x = st[pa[y]])

 $REP(x\,,\ 1\,,\ nx)\ g\,[\,b\,]\,[\,x\,]\,.w\,=\,g\,[\,x\,]\,[\,b\,]\,.w\,=\,0\,;$

f.pb(x), f.pb(y = st[match[x]]), $q_push(y)$;

```
\begin{array}{l} \label{eq:condition} \text{fill} \left( \text{ALL} \left( \text{flo\_from} \left[ b \right] \right) \,, \,\, 0 \right); \\ \text{for} \,\, \left( \text{int} \,\, xs \,: \,\, \text{flo} \left[ b \right] \right) \,\, \left\{ \end{array}
          g[b][x] = g[xs][x], g[x][b] = g[x][xs];
                              1, n)
                 if (flo_from[xs][x]) flo_from[b][x] = xs;
     set_slk(b);
void expand_blossom(int b) {
     for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
     for (int x: split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
           slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
      for (int x : flo[b])
            if (x = xr) S[x] = 1, pa[x] = pa[b];
            else S[x] = -1, set\_slk(x);
      \operatorname{st}[b] = 0;
bool on_found_edge(const edge &e) {
      \begin{array}{lll} & \text{if (int } u = st \, [e \cdot u] \;, \; v = st \, [e \cdot v] \;; \; S[v] = -1) \; \{ \\ & \text{int } nu = st \, [match[v]] \;; \; pa[v] = e \cdot u \;; \; S[v] = 1 \;; \\ & slk \, [v] = slk \, [nu] = S[nu] = 0 \;; \; q\_push(nu) \;; \\ \end{array} 
      else if (S[v] = 0)
            if (int o = lca(u, v)) add_blossom(u, o, v);
           else return augment(u, v), augment(v, u), true;
     return false;
bool matching() { fill(ALL(S), -1), fill(ALL(slk), 0);
      q = queue < int > ();
     \begin{array}{l} \text{REP}(x, 1, nx) & \text{if } (st[x] = x \& x : match[x]) \\ pa[x] = S[x] = 0, \text{ q_push}(x); \end{array}
       if (q.empty()) return false;
      for (;;)
            while (SZ(q)) {
                 int u = q.front(); q.pop();
                  if (S[st[u]] = 1) continue;
                REP(v, 1, n)
                       if (g[u][v].w > 0 & st[u] != st[v]) {
                             if (E(g[u][v]) != 0)
                                  update\_slk(u, st[v], slk[st[v]]);
                                            (on_found_edge(g[u][v])) return true;
            int d = INF;
          REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
           REP(x,\ 1,\ nx)
                 if (int
                      s = slk[x]; st[x] == x && s && s[x] <= 0)

d = min(d, E(g[s][x]) / (S[x] + 2));
           REP(u, 1, n)
                 if (S[st[u]] = 1) lab[u] += d;
                 else if (S[st[u]] = 0) {
   if (lab[u] \le d) return false;
                       lab[u] -= d;
          REP(x, 1, nx)
                 if (int s = slk[x]; st[x] == x &&
                       \begin{array}{c} s \&\& st[s] \mathrel{!=} x \&\& E(g[s][x]) == 0) \\ if (on\_found\_edge(g[s][x])) \ return \ true; \end{array}
           REP(b\,,\ n\,+\,1\,,\ nx\,)
                  if (st[b] = b \&\& S[b] = 1 \&\& lab[b] = 0)
                      expand_blossom(b);
     return false;
pair < ll, int > solve() {
      fill (ALL(match), 0);
      \begin{array}{lll} \text{REP}(\mathbf{u}, \ \mathbf{0}, \ \mathbf{n}) & \text{st} \ [\mathbf{u}] = \mathbf{u}, \ \text{flo} \ [\mathbf{u}]. \ \text{clear} \ (); \\ & \text{int} \ \mathbf{w} \_ \text{max} = \ \mathbf{0}; \\ \end{array} 
      \begin{array}{lll} \text{REP}(u, 1, n) & \text{REP}(v, 1, n) & \{ & \\ \text{flo\_from} \left[u\right] \left[v\right] & = \left(u = v ? u : 0\right); \end{array} 
           w_m = max(w_m x, g[u][v].w);
      fill (ALL(lab), w_max);
     \begin{tabular}{ll} \beg
```

return 0;

bool bfs() {

 $fill_n(dis, n + 3, -1);$

```
\begin{array}{l} \text{queue} < \text{int} > \text{q}; \\ \text{q.push}(\text{s}), & \text{dis}[\text{s}] = 0; \end{array}
      while (matching()) ++n_matches;
     REP(u, 1, n) if (match[u] \&\& match[u] < u)
        tot_weight += g[u][match[u]].w;
                                                                                      while (!q.empty()) {
     return make_pair(tot_weight, n_matches);
                                                                                         int u = q. front();
                                                                                         q.pop();
  void add_edge(int u, int v, int w)
{ g[u][v].w = g[v][u].w = w; }
                                                                                         for (edge &e : G[u])
if (!~dis[e.to] && e.flow != e.cap)
                                                                                              q.push(e.to), dis[e.to] = dis[u] + 1;
4.6 SW-mincut [90bfe6]
                                                                                      return dis[t] != -1;
int maxflow(int _s, int _t) {
                                                                                      s = \_s, t = \_t;
                                                                                      int \overline{flow} = 0, df;
  \begin{array}{ll} \text{int vst} \left[ \text{MXN} \right], \text{ edge} \left[ \text{MXN} \right] \left[ \text{MXN} \right], \text{ wei} \left[ \text{MXN} \right]; \\ \text{void init} \left( \text{int n} \right) \end{array} \right\}
                                                                                      while (bfs()) {
                                                                                         fill_n(cur, n + 3, 0);
while ((df = dfs(s, INF))) flow += df;
     REP fill_n(edge[i], n, 0);
   void addEdge(int u, int v, int w){
                                                                                      return flow;
     edge[u][v] += w; edge[v][u] += w;
                                                                                   bool solve() {
   int search(int &s, int &t, int n){
                                                                                      int sum = 0;
     fill_n(vst, n, 0), fill_n(wei, n, 0);
                                                                                      for (int i = 0; i < n; ++i)
     s = t = -1;
                                                                                         if (cnt[i] > 0)
     int mx, cur;
                                                                                      add_edge(n + 1, i, cnt[i]), sum += cnt[i];
else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int j = 0; j < n; +++j) {
       mx = -1, cur = 0;

mx = -1, cur = 0;

mx = -1, mx = wei[i];
                                                                                      for (int i = 0; i < n; ++i)
        vst[cur] = 1, wei[cur] = -1;
                                                                                             (\,\mathrm{cnt}\,[\,i\,]\,>\,0)
        s = t; t = cur;
                                                                                         G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)
        G[i].pop\_back(), G[n + 2].pop\_back();
     return mx;
                                                                                      return sum != -1;
   int solve(int n) {
                                                                                   int solve(int _s, int
                                                                                                                  t) {
     int res = INF;
                                                                                      add_edge(_t, _s, INF);
if (!solve()) return -1; // invalid flow
     for (int x, y; n > 1; n--){
         \begin{array}{l} res = min(res, search(x, y, n)); \\ REP \ edge[i][x] = (edge[x][i] += edge[y][i]); \\ \end{array} 
                                                                                      int x = G[\underline{t}] . back() . flow;
                                                                                      return G[_t].pop_back(), G[_s].pop_back(), x;
        REP {
          edge[y][i] = edge[n - 1][i];
edge[i][y] = edge[i][n - 1];
                                                                                };
          // edge[y][y] = 0; 
                                                                                4.8 Gomory Hu tree* [5f2460]
                                                                                MaxFlow Dinic;
     return res;
                                                                                int g[MAXN];
  }
} sw;
                                                                                void GomoryHu(int n) { // 0-base
                                                                                   fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
4.7
       BoundedFlow*(Dinic*) [4ae8ab]
                                                                                      Dinic.reset();
{\tt struct} \;\; {\tt BoundedFlow} \; \left\{ \;\; // \;\; {\tt 0-base} \right.
                                                                                      add\_edge(\,i\;,\;g\,[\,i\,]\;,\;Dinic.maxflow(\,i\;,\;g\,[\,i\,]\,)\,)\,;
   struct edge {
                                                                                      for (int j = i + 1; j \le n; ++j)

if (g[j] = g[i] \&\& \sim Dinic.dis[j])
     int to, cap, flow, rev;
                                                                                           g[j] = i;
   vector<edge> G[N];
   int n, s, t, dis [N], cur [N], cnt [N];
                                                                                }
   void init(int _n) {
     n = _n;
for (int i = 0; i < n + 2; ++i)
                                                                                       Minimum Cost Circulation* [cb40c6]
        G[i].clear(), cnt[i] = 0;
                                                                                struct MinCostCirculation { // 0-base
                                                                                   {\tt struct} \  \, {\tt Edge} \  \, \{
   void add_edge(int u, int v, int lcap, int rcap) {
                                                                                      ll from, to, cap, fcap, flow, cost, rev;
     } *past[N];
                                                                                    vector<Edge> G[N];
                                                                                   ll dis [N], inq [N], n;
void BellmanFord(int s) {
  void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, SZ(G[v])});
  G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
                                                                                      fill_n(dis, n, INF), fill_n(inq, n, 0);
                                                                                      queue<int> q;
                                                                                      auto relax = [&](int u, ll d, Edge *e) {
                                                                                         if (dis[u] > d) {
   int dfs(int u, int cap) {
  if (u = t || !cap) return cap;
                                                                                           \begin{array}{l} \text{dis}\,[u] = d, \;\; \text{past}\,[u] = e; \\ \text{if} \;\; (!\, \text{inq}\,[u]) \;\; \text{inq}\,[u] = 1, \;\; \text{q.push}(u); \end{array}
     for (int &i = cur[u]; i < SZ(G[u]); ++i) {
                                                                                        }
        edge \&e = G[u][i];
        if (dis[e.to] == dis[u] + 1 && e.cap!= e.flow) {
                                                                                      relax(s, 0, 0);
           int df = dfs(e.to, min(e.cap - e.flow, cap));
                                                                                      while (!q.empty()) {
           if (df) {
                                                                                         int u = q.front();
              e.flow += df, G[e.to][e.rev].flow -= df;
                                                                                         q.pop(), inq[u] = 0;
              return df;
                                                                                         for (auto &e : G[u])
          }
                                                                                           if(e.cap > e.flow)
                                                                                              relax\left(\begin{smallmatrix}e.to\end{smallmatrix}, \; dis\left[\begin{matrix}u\end{smallmatrix}\right] \; + \; e.cost\;, \; \&e\right);
        }
     dis[u] = -1;
```

void try_edge(Edge &cur) {

if (dis[cur.from] + cur.cost < 0) {

BellmanFord(cur.to);

if (cur.cap > cur.flow) return ++cur.cap, void();

4.10 Flow Models

- $\bullet \quad {\rm Maximum/Minimum\,flow\,with\,lower\,bound/Circulation\,problem}$
 - 1. Construct supersource S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - − To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite $\operatorname{graph}(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- $\bullet \quad {\rm Minimum} \, {\rm cost} \, {\rm cyclic} \, {\rm flow}$
 - 1. Consruct super source S and $\sinh T$
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - $4. \ \ \text{For each vertex} \ v \ \text{with} \ d(v) > 0, \\ \text{connect} \ S \rightarrow v \ \text{with} \ (cost, cap) = (0, d(v))$
 - 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow

- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$
 $p_{u} \ge 0$

5 String

5.1 KMP [9e1cd1]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
}
for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
}
return ans;
}</pre>
```

5.2 Z-value* [e2dc6f]

5.3 Manacher* [bfe74e]

5.4 SAIS* [e9a275]

```
auto sais (const auto &s) {
  const int n = SZ(s), z = ranges::max(s) + 1;
  if (n = 1) return vector\{0\};
vector<int> c(z); for (int x : s) ++c[x];
  partial_sum(ALL(c), begin(c));
  vector < int > sa(n); auto I = views :: iota(0, n);
  vector<br/>bool> t(n, true);
  for (int i = n - 2; i \ge 0; --i)
    t[i] = (
         s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
  auto is_lms = views::filter([&t](int x) {
    return x && t[x] && !t[x - 1];
  });
  auto induce = [&] {
    for (auto x = c; int y : sa)

if (y--) if (!t[y]) sa[x[s[y]-1]++] = y;

for (auto x = c; int y : sa | views::reverse)
       if (y--) if (t[y]) sa[--x[s[y]]] = y;
  vector < int > lms, q(n); lms.reserve(n);
  for (auto x = c; int i : I | is_lms)
    q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
  induce(); vector<int> ns(SZ(lms));
```

for (int j = -1, nz = 0; int $i : sa \mid is_lms$) {

```
if (j >= 0) {
        int len = min({n - i, n - j, lms[q[i] + 1] - i});
        ns[q[i]] = nz += lexicographical_compare(
             begin(s) + j, begin(s) + j + len
             begin(s) + i, begin(s) + i + len);
     j \; = \; i \; ;
   fill(ALL(sa), 0); auto nsa = sais(ns);
   for (auto x = c; int y : nsa | views::reverse)
     y = lms[y], sa[--x[s[y]]] = y;
   return induce(), sa;
// sa[i]: sa[i]-th suffix
       is the i-th lexicographically smallest suffix.
   hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
  int n; vector<int> sa, hi, ra;
   Suffix
     (const auto &_s, int _n) : n(_n), hi(_n), ra(_n) { vector < int > s(_n + 1); // s[_n] = 0; copy_n(_s, n, begin(s)); // _s shouldn't contain 0
     sa = sais(s); sa.erase(sa.begin());
     for (int i = 0; i < n; ++i) ra[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {
  if (!ra[i]) { h = 0; continue; }
  for (int j = sa[ra[i] - 1]; max
        (i, j) + h < n & & s[i + h] == s[j + h];) ++h;

hi[ra[i]] = h ? h-- : 0;
  }
};
```

5.5 Aho-Corasick Automatan* [91c6c0]

```
struct AC_Automatan {
   int nx[len][sigma], fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
   int newnode() {
      fill_n(nx[top], sigma, -1);
      return top++;
   void init() { top = 1, newnode(); }
   int input(string &s) {
      int X = 1;
      \begin{array}{lll} & \text{for } (c\text{ har } c:s) \ \{ & \text{if } (!\!\sim\!\!\operatorname{nx}[X][c-'A']) \ \operatorname{nx}[X][c-'A'] = \operatorname{newnode}() \,; \\ & X = \operatorname{nx}[X][c-'A'] \,; \end{array}
      return X; // return the end node of string
   void make_fl() {
      queue\langle int \rangle q;
q.push(1), fl[1] = 0;
       for (int t = 0; !q.empty(); ) {
          int R = q.front();
          q.pop(), ord[t++] = R;
          for (int i = 0; i < sigma; ++i)
if (~nx[R][i]) {
                int X = rnx[R][i] = nx[R][i], Z = fl[R];
                 for (; Z \&\& !\sim nx[Z][i];) Z = fl[Z];

fl[X] = Z ? nx[Z][i] : 1, q.push(X);
             else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
      }
   void solve() {
      \begin{array}{lll} & \text{for (int i = top - 2; i > 0; --i)} \\ & \text{cnt[fl[ord[i]]] += cnt[ord[i]];} \end{array}
} ac;
```

5.6 Smallest Rotation [e74dc0]

```
\begin{array}{l} string \ mcp(string \ s) \ \{\\ int \ n = SZ(s) \,, \ i = 0 \,, \ j = 1;\\ s += s \,;\\ while \ (i < n \&\& \ j < n) \ \{\\ int \ k = 0 \,;\\ while \ (k < n \&\& \ s[i + k] \Longrightarrow s[j + k]) \ +\!\!\!+k;\\ if \ (s[i + k] <= s[j + k]) \ j +\!\!\!= k + 1;\\ else \ i +\!\!\!= k + 1 \,;\\ if \ (i \Longrightarrow j) +\!\!\!+j \,;\\ \}\\ int \ ans = i < n \ ? \ i \ : j \,;\\ return \ s.substr(ans, n) \,;\\ \} \end{array}
```

5.7 De Bruijn sequence* [f601c2]

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N
  void dfs(int *out, int t, int p, int &ptr) {
    if (ptr >= L) return;
    if (t > N) {
        if (N % p) return;
        for (int i = 1; i <= p && ptr < L; ++i)
            out[ptr++] = buf[i];
    } else {
        buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
        for (int j = buf[t - p] + 1; j < C; ++j)
            buf[t] = j, dfs(out, t + 1, t, ptr);
        }
    }
    void solve(int _c, int _n, int _k, int *out) {
        int p = 0;
        C = _c, N = _n, K = _k, L = N + K - 1;
        dfs(out, 1, 1, p);
        if (p < L) fill(out + p, out + L, 0);
    }
} dbs;</pre>
```

```
5.8 Extended SAM* [58fa19]
struct exSAM { int len [N * 2], link [N * 2]; // maxlength, suflink int next [N * 2] [CNUM], tot; // [0, tot), root = 0 int lenSorted [N * 2]; // topo. order
   int cnt[N * 2]; // occurence
   int newnode()
     fill_n (next[tot], CNUM, 0);
     len[tot] = cnt[tot] = link[tot] = 0;
     return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
  int cur = next[last][c];
     len[cur] = len[last] + 1;
int p = link[last];
     while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
     int q = next[p][c];
     if (len
          [p] + 1 \Longrightarrow len[q] return link[cur] = q, cur;
     int clone = newnode();
     for (int i = 0; i < CNUM; ++i)
       next
             [i] = len[next[q][i]] ? next[q][i] : 0;
     len[clone] = len[p] + 1;
     while (p != -1 \&\& next[p][c] == q)
     next[p][c] = clone, p = link[p];
link[link[cur] = clone] = link[q];
     link[q] = clone;
     return cur;
   void insert(const string &s) {
     int cur = 0:
     for (auto ch : s) {
       int & mxt = next[cur][int(ch - 'a')];
        if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
  void build() {
     queue < int > q;
     q.push(0);
     while (!q.empty()) {
       int cur = q.front();
        q.pop();
        for (int i = 0; i < CNUM; ++i)
          if \(\text[cur][i])
            q.push(insertSAM(cur, i));
     vector<int> lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];
     partial_sum(ALL(lc), lc.begin());
     for (int i
          = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  void solve() {
     for (int i = tot - 2; i >= 0; --i)
```

cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];

} else {

l = cntr - nu + 1;

```
\begin{array}{l} k1 \, = \, {\rm get\_z} \big( {\rm z3} \, , \, \, nu \, + \, 1 \, + \, nv \, - \, 1 \, - \, (\, {\rm cntr} \, \, - \, nu) \, \big) \, ; \\ k2 \, = \, {\rm get\_z} \big( {\rm z4} \, , \, \, (\, {\rm cntr} \, \, - \, nu) \, + \, 1 \big) \, ; \end{array}
| };
 5.9 PalTree* [675736]
                                                                       if (k1 + k2 >= 1)
 struct palindromic_tree {
                                                                         add_rep(cntr < nu, cntr, l, k1, k2);
   struct node {
     int next[26], fail, len;
                                                                  int cnt, num; // cnt: appear times, num: number of
    // pal. suf.
                                                                       Math
     node(int l = 0): fail(0), len(1), cnt(0), num(0) {
                                                                  6.1 ax+by=gcd(only exgcd *) [5fef50]
       for (int i = 0; i < 26; ++i) next[i] = 0;
                                                                   pll exgcd(ll a, ll b)
                                                                     if (b = 0) return pll(1, 0);
   vector<node> St;
                                                                     ll p = a / b;
   vector<char> s;
                                                                     pll q = exgcd(b, a \% b);
   int last, n;
                                                                     return pll(q.Y, q.X - q.Y * p);
   \begin{array}{lll} palindromic\_tree() : St(2), \ last(1), \ n(0) \ \{ \\ St[0].fail = 1, \ St[1].len = -1, \ s.pb(-1); \end{array}
                                                                   /* ax+by=res, let x be minimum non-negative
                                                                  g, p = gcd(a, b), exgcd(a, b) * res / g
if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
   inline void clear() {
     St.clear(), s.clear(), last = 1, n = 0; St.pb(0), St.pb(-1);
                                                                  else: t = -(p.X / (b / g))

p += (b / g, -a / g) * t */
     St[0].fail = 1, s.pb(-1);
                                                                  6.2 Floor and Ceil [1ffa73]
   inline int get_fail(int x) {
     while (s[\tilde{n} - St[x].len - 1] != s[n])
                                                                  x = St[x].fail;
     return x;
                                                                  { return a / b + (a \% b && (a < 0) ^{\circ} (b > 0)); }
   inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
                                                                  6.3 Floor Enumeration [67ad61]
     int cur = get_fail(last);
                                                                    / enumerating x = floor(n / i), [1,
     if (!St[cur].next[c]) {
                                                                   for (int l = 1, r; l \le n; l = r + 1) {
        int now = SZ(St);
                                                                    int x = n / 1;
        St.pb(St[cur].len + 2);
                                                                    r = n / x;
                                                                  }
       St [now]. fail =
          St[get_fail(St[cur].fail)].next[c];
                                                                  6.4 Mod Min [038fef]
        St[cur].next[c] = now;
       St[now].num = St[St[now].fail].num + 1;
                                                                    / \min\{k \mid l \le ((ak) \mod m) \le r\}, \text{ no solution } -> -1
                                                                  if (a == 0) return 1 ? -1 : 0;

if (ll k = (l + a - 1) / a; k * a <= r)
     last = St[cur].next[c], ++St[last].cnt;
   inline void count() { // counting cnt
                                                                       return k;
     auto i = St.rbegin();
                                                                     ll b = m / a, c = m \% a;
     for (; i != St.rend(); ++i) {
                                                                     St[i->fail].cnt += i->cnt;
                                                                       return (1 + y * c + a - 1) / a + y * b;
                                                                     return -1:
                                                                 }
   inline int size() { // The number of diff. pal.
                                                                  6.5 Gaussian integer gcd [4fcbff]
     return SZ(St) - 2;
                                                                  cpx \ gaussian\_gcd(cpx \ a\,, \ cpx \ b) \ \{
                                                                  #define rnd
                                                                       (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
 5.10 Main Lorentz [eaf279]
                                                                     ll d = a.imag() * b.real() + a.imag() * b.imag();
ll d = a.imag() * b.real() - a.real() * b.imag();
 vector<pair<int, int>>> rep[kN]; // 0-base [l, r]
                                                                     ll r = b.real() * b.real() + b.imag() * b.imag();
 void main_lorentz(const string &s, int sft = 0) {
                                                                     if (c % r = 0 & d % r = 0) return b;
   const int n = s.size();
   if (n = 1) return;
                                                                     return gaussian_gcd
   (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
   const string u = s.substr(0, nu), v = s.substr(nu),
          ru(u.rbegin
                                                                  6.6 GaussElimination [c016c9]
              (), u.rend()), rv(v.rbegin(), v.rend());
   main_lorentz(u, sft), main_lorentz(v, sft + nu);
                                                                   void GAS(V<V<double>>&vc) {
   const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u)
                                                                       int len = vc.size();
   for (int i = 0; i < len; ++i)
                                                                            int idx = find_if(vc.begin()+i,vc.end(),[\&](
     return
                                                                                auto&v) {return v[i] != 0;} ) - vc.begin();
           (0 \le i \text{ and } i < (int)z.size()) ? z[i] : 0; };
                                                                            if ( idx == len ) continue;
   auto add_rep
                                                                            if( i != idx ) swap( vc[idx], vc[i] );
double pivot = vc[i][i];
        =\,[\,\&\,](\,bool\ left\;,\;int\;\,c\,,\;int\;\,l\,,\;int\;\,k1\,,\;int\;\,k2)\;\,\{
                                                                           const
           int L = \max(1, 1 - k2), R = \min(1 - left, k1);
     if (L > R) return;
     if (left)
           rep[1].emplace\_back(sft + c - R, sft + c - L);
     else rep[l].emplace_back
                                                                                     double mul = vc[j][i]/vc[i][i];
transform( vc[j].begin(), vc[j].end
          (sft + c - R - l + 1, sft + c - L - l + 1);
                                                                                          ()\;,\;\;vc\,[\,i\,]\;.\;begin\,()\;,\;\;vc\,[\,j\,]\;.\;begin\,()\;,
   for (int cntr = 0; cntr < n; cntr++) {
                                                                                              [&](auto &a, auto &b) {
     int 1, k1, k2;
                                                                                              return a-b*mul;
     if (cntr < nu) {
                                                                                              });
        l = nu - cntr;
                                                                                }
       k1 = get_z(z1, nu - cntr);

k2 = get_z(z2, nv + 1 + cntr);
                                                                           }
                                                                       }
```

};

6.7 Miller Rabin* [14b81a]

```
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool Miller_Rabin(ll a, ll n) {
  if ((a = a \% n) = 0) return 1;
  if (n \% 2 = 0) return n = 2;
  ll tmp = (n - 1) / ((n - 1) & (1 - n));
ll t = \underline{lg(((n - 1) & (1 - n)))}, x = 1;
  for (; tmp; tmp >>= 1, a = mul(a, a, n))
    if (tmp \& 1) x = mul(x, a, n);
  if (x = 1 | | x = n - 1) return 1;
  while (--t)
    if ((x = mul(x, x, n)) = n - 1) return 1;
```

6.8 Simultaneous Equations [21b2e1]

```
struct matrix { //m variables, n equations
   int n, m;
   \label{eq:maxn} \textit{fraction} \ M[\textit{MAXN}] \left[ \textit{MAXN} + \ 1 \right], \ \textit{sol} \left[ \textit{MAXN} \right];
   int solve() { //-1: inconsistent, >= 0: rank for (int i = 0; i < n; ++i) {
        int piv = 0;
         while (piv < m && !M[i][piv].n) ++piv;
         if (piv == m) continue;
         for (int j = 0; j < n; +++j) {
           if (i == j) continue;
            \begin{array}{ll} \text{fraction tmp} = \text{-M[j][piv]} \ / \ \text{M[i][piv];} \\ \text{for (int } k = 0; \ k <= \\ \end{array} 
                  m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
        }
      int rank = 0;
      for (int i = 0; i < n; ++i) {
        int piv = 0;
         while (piv < m && !M[i][piv].n) ++piv;
        if (piv = m & M[i][m].n) return -1;
        {\color{red} {\rm else}} if (piv
               < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
      return rank;
  }
};
```

6.9 Pollard Rho* [fff0fc]

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n = 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
      = 0) return PollardRho(n / 2), ++cnt[2], void();
  11 x = 2, y = 2, d = 1, p = 1;
 #define f(x, n, p) ((mul(x, x, n) + p) % n) while (true) {
    if (d!= n & d!= 1) {
      PollardRho(n / d);
      PollardRho(d);
      return;
    if (d = n) + p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);

d = gcd(abs(x - y), n);
```

6.10 Simplex Algorithm [40618e]

```
const double eps = 1E-10;
double a [MAXN] [MAXM] , b [MAXN] , c [MAXM] ; double d [MAXN] [MAXM] , x [MAXM] ; int ix [MAXN + MAXM] ; // !!! array all indexed from 0
// \max\{cx\} \text{ subject to } \{Ax \le b, x > = 0\}
  n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  fill_n(d[n], m + 1, 0);
  fill_n(d[n+1], m+1, 0);
  iota(ix, ix + n + m, 0);
```

```
int r = n, s = m - 1;
for (int i = 0; i < n; ++i) {
  for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j]; d[i][m - 1] = 1; d[i][m] = b[i];
  if (d[r][m] > d[i][m]) r = i;
copy_n(c, m - 1, d[n]);

d[n + 1][m - 1] = -1;
for (double dd;; ) {
  if (r < n) {
    swap(ix[s], ix[r+m])
    for (int i = 0; i \le n + 1; ++i) if (i != r) {
  for (int j = 0; j \le m; ++j) if (j != s)
  d[i][j] += d[r][<math>j] * d[i][s];
  d[i][s] *= d[r][s];
    }
  }
  r = s = -1;
  if (s < 0) break;
  for (int \ i = 0; \ i < n; ++i) \ if \ (d[i][s] < -eps) {
    if^{(r < 0)}
         (dd = d[r][m]
               / d[r][s] - d[i][m] / d[i][s]) < -eps | |
         (dd < eps & ix[r + m] > ix[i + m]))
       r = i;
  if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not executable
double ans = 0;
fill_n(x, m, 0);
\quad \text{for (int } i = m; \ i <
  n+m; ++i) { // the missing enumerated x[i] = 0 if (ix[i] < m - 1){
    ans += d[i - m][m] * c[ix[i]];
    x[ix[i]] = d[i-m][m];
}
return ans;
```

6.10.1 Construction

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c$

 $\overline{\mathbf{x}}$ and $\overline{\mathbf{y}}$ are optimalified only if for all $i \in [1, n]$, either $\overline{x}_i = 0$ or $\sum_{j=1}^m A_{ji} \overline{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij} \bar{x}_j = b_j$ holds.

```
1. In case of minimization, let c'_i = -c_i
```

```
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
```

- $3. \sum_{1 \le i \le n}^{-} A_{ji} x_i = b_j$
- $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
- $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.11 chineseRemainder [fe9f25]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
   11 g = gcd(m1, m2);
   if ((x2 - x1) % g) return -1; // no sol
   m1 /= g; m2 /= g;
   \begin{array}{l} {\rm pll} \ p = {\rm exgcd} \, (m1, \ m2) \, ; \\ {\rm ll} \ {\rm lcm} = {\rm m1} \ ^* \ {\rm m2} \ ^* \, {\rm g} \, ; \\ {\rm ll} \ {\rm res} = p. \, {\rm first} \ ^* \, ({\rm x2} \ - \, {\rm x1}) \ ^* \, {\rm m1} + {\rm x1} \, ; \end{array}
   // be careful with overflow
   return (res % lcm + lcm) % lcm;
```

6.12 Factorial without prime factor* [dcffcb]

```
O(p^k + \log^2 n), pk = p^k
ll prod [MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
 \operatorname{prod}[0] = 1;
  for (int i = 1; i \le pk; ++i)
    if (i \% p) prod[i] = prod[i - 1] * i \% pk;
```

```
else prod[i] = prod[i - 1];
   11 \text{ rt} = 1;
  for (; n; n /= p) {
   rt = rt * mpow(prod[pk], n / pk, pk) % pk;
   rt = rt * prod[n % pk] % pk;
} // (n! without factor p) % p^k
```

6.13 Discrete Log* [ba4ac0]

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered\_map < int , int > p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {
     p[y] = i;
y = 1LL * y * x % m;
b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;
      if \ (p.find(s) \mathrel{!=} p.end()) \ return \ i \ + \ kStep \ - \ p[\, s \,]; \\
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
     if (s == y) return i;
s = 1LL * s * x % m;
  if (s = y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
```

6.14 Berlekamp Massey [9380b8]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
   vector(I) Bertekamphassey (const vector(I) & output vector(I) d(SZ(output) + 1), me, he; for (int f = 0, i = 1; i <= SZ(output); ++i) { for (int j = 0; j < SZ(me); ++j) d[i] += output[i - j - 2] * me[j]; if ((d[i] -= output[i - 1]) == 0) continue; if (me approximation) [
        if (me.empty()) {
  me.resize(f = i);
             continue;
       vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
for (T x : he) o.pb(x * k);
        o.resize (\max(SZ(o), SZ(me)));
for (\text{int } j = 0; j < SZ(me); +++j) o[j] += me[j];
        if (i - f + SZ(he)) >= SZ(me) he = me, f = i;
        me = o:
    return me;
```

6.15 Primes

```
12721 13331 14341 75577 123457 222557
  556679 999983 1097774749 1076767633 100102021
 999997771 1001010013 1000512343 987654361 999991231
   999888733 98789101 987777733 999991921 1010101333
   1010102101 1000000000039 100000000000037
   2305843009213693951 4611686018427387847
   9223372036854775783 18446744073709551557 */
```

6.16 Theorem

• Cramer's rule

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i), L_{ij} = -c$ where cisthenumber of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence $d_1, d_2, ..., d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!} \text{ spanning trees.}$ Let $T_{n,k}$ be the number of labeled forests on n vertices with k composite labeled forests on labeled fo
 - nents, such that vertex 1,2,...,k belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős-Gallaitheorem

As equence of nonnegative integers $d_1 \ge \cdots \ge d_n$ can be represented as the ${\it degree sequence of a finite simple graph on } n {\it vertices} if and only if d_1 + \dots + d_n$

degreesequenceofafinitesimplegraphon
$$n$$
verticesifandonlyif $d_1+\cdots+$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$ holds for every $1 \leq k \leq n$. Gale–Byser theorem

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \ge \cdots \ge a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^{n} a_i = \sum_{i=1}^{n} b_i$ and $\sum_{i=1}^{k} a_i \leq \sum_{i=1}^{n} \min(b_i, k)$ holds for every $1 \le k \le n$.

Fulkerson-Chen-Ansteetheorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1) + \sum_{i=k+1}^n \min(b_i, k) \operatorname{holds} \operatorname{for} \operatorname{every} 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have A =#{lattice points in the interior}+ $\frac{\#\{\text{lattice points on the boundary}\}}{2}-1$.

- Möbius inversion formula

 - $\begin{array}{ll} & f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d}) \\ & f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d) \end{array}$
- - A portion of a sphere cut off by a plane.
 - -r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : $\arcsin(a/r)$.
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \sin \theta)$ $\cos\theta)^2/3$.
 - Area = $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2 (1 \cos\theta)$.
- Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n) = 0$.
 - Lagrangian function $\mathcal{L}(x_1, \ldots, x_n, \lambda_1, \ldots, \lambda_k) = f(x_1, \ldots, x_n)$ $\sum_{i=1}^{k} \lambda_i g_i(x_1, ..., x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
 - Line 1: $v_1 = p_1 + t_1 d_1$
 - Line 2: $\boldsymbol{v}_2 = \boldsymbol{p}_2 + t_2 \boldsymbol{d}_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
 - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$
 - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$

 - $egin{array}{l} n_2\!=\!a_2\! imes\!n_1 \\ -c_1\!=\!p_1\!+\!rac{(p_2\!-\!p_1)\!\cdot\!n_2}{d_1\!\cdot\!n_2}d_1 \\ -c_2\!=\!p_2\!+\!rac{(p_1\!-\!p_2)\!\cdot\!n_1}{d_2\!\cdot\!n_1}d_2 \end{array}$
- Derivatives/Integrals

Integration by parts: $\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$ $\frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^{2}}} \begin{vmatrix} \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^{2}}} \\ \int \tan x = 1 + \tan^{2}x \end{vmatrix} \int \tan x = -\frac{\ln|\cos x|}{a}$ $\int e^{-x^{2}} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \begin{vmatrix} \int xe^{ax}dx = \frac{e^{ax}}{a^{2}}(ax-1) \end{vmatrix}$ $\int \sqrt{a^2 + x^2} = \frac{1}{2} \left(x \sqrt{a^2 + x^2} + a^2 \sinh(x/a) \right)$

• Spherical Coordinate

 $(x,y,z) = (r\sin\theta\cos\phi, r\sin\theta\sin\phi, r\cos\theta)$

$$(r,\!\theta,\!\phi)\!=\!(\sqrt{x^2\!+\!y^2\!+\!z^2},\! \mathrm{acos}(z/\sqrt{x^2\!+\!y^2\!+\!z^2}),\! \mathrm{atan2}(y,\!x))$$

• Rotation Matrix

$$M(\theta) = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix}, R_x(\theta_x) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \cos\theta_x & -\sin\theta_x \\ 0 & \sin\theta & \cos\theta \end{bmatrix}$$

6.17 Estimation

n | 2345 6 7 8 9 20 30 40 50 100 p(n) = 23571115223062756044e42e52e8 $n \hspace{0.2cm} | \hspace{0.02cm} 100\hspace{0.02cm} 1 \hspace{0.02cm} e3\hspace{0.02cm} 1 \hspace{0.02cm} e6\hspace{0.02cm} 1 \hspace{0.02cm} e9\hspace{0.02cm} 1 \hspace{0.02cm} e12\hspace{0.02cm} 1 \hspace{0.02cm} e15\hspace{0.02cm} \hspace{0.02cm} 1 \hspace{0.02cm} e18\hspace{0.02cm}$ d(i) 12 32 2401344672026880103680 $n \mid 123456789$ 10 11 12 13 14 15 $\binom{2n}{n}$ 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8 $n \mid 23456789$ 10 11 12 13 B_n 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7

6.18 Euclidean Algorithms

- $m = |\frac{an+b}{a}|$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) = & \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor \\ = & \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \\ g(a,b,c,n) = & \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ = & \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ = & \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases} \\ h(a,b,c,n) = & \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2} \\ = & \begin{cases} \lfloor \frac{a}{c} \rfloor^{2} \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^{2} \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.19 General Purpose Numbers

• Bernoullinumbers

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 $S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$ $x^n = \sum_{i=0}^{n} S(n,i)(x)_i$ • Pentagonal number theorem

• Pentagonal number theorem
$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
• Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. kj:ss.t. $\pi(j) > \pi(j+1), k+1j$:ss.t. $\pi(j) \ge j$, kj:ss.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.20 Tips for Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i \ge 0} a_i x^i$
 - $-A(rx) \Rightarrow r^n a_n$

 - $-A(x)+B(x) \Rightarrow a_n+b_n$ -A(x)B(x)\Rightarrow\sum_{i=0}^n a_i b_{n-i}
 - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $-xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i > 0} \frac{a_i}{i!} x_i$

- $-A(x)+B(x) \Rightarrow a_n+b_n$ $-A^{(k)}(x) \Rightarrow a_{n+k}$ $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$ $-A(x)^{k} \Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n} {n \choose i_{1},i_{2},\dots,i_{k}} a_{i_{1}} a_{i_{2}} \dots a_{i_{k}}$ $-xA(x) \Rightarrow na_n$
- Special Generating Function
 - $(1+x)^n = \sum_{i \ge 0} {n \choose i} x^i$ $\frac{1}{(1-x)^n} = \sum_{i \ge 0} {i \choose n-1} x^i$

Polynomial

Fast Fourier Transform [9fec80]

```
const int maxn = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = acos(-1);
cplx omega[maxn + 1];
\begin{array}{c} \text{void prefft() } \{\\ \text{ for (int } i = 0; i <= \max \\ \end{array}
          ; ++i) omega[i] = exp(i * 2 * pi / maxn * I);
void bin(vector<cplx> &a, int n) {
     int lg;
    \begin{array}{ll} \text{for (lg = 0; (1 << lg) < n; ++lg); --lg;} \\ \text{vector} < \text{cplx} > \text{tmp(n);} \end{array}
     for (int i = 0; i < n; ++i) {
         int to = 0;
         for (int j = 0; (1 << j) <
n; ++j) to |= (((i >> j) & 1) << (lg - j));
         tmp[to] = a[i];
     for (int i = 0; i < n; ++i) a[i] = tmp[i];
}
void fft(vector<cplx> &a, int n) {
     bin(a, n);
     for (int step = 2; step \leq n; step \leq 1) {
         int to = step \gg 1;
         for (int i = 0; i < n; i += step) {
              for (int k = 0; k < to; ++k) {
                   cplx x = a[i
                      + to + k | * omega [maxn / step * k];
                   a[i + to + k] = a[i + k] - x;
                   a[i + k] += x;
              }
         }
void ifft (vector<cplx> &a, int n) {
     fft(a, n);
     reverse(a.begin() + 1, a.end());
     for (int i = 0; i < n; ++i) a[i] /= n;
vector<int> multiply(const vector<
     int>&a, const vector<int> &b, bool trim = false) {
     int d = 1;
     while
          (d < max(a.size(), b.size())) d <<= 1; d <<= 1;
     vector < cplx > pa(d), pb(d);
     for (int i
          = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
     for (int i
           = 0; i < b.size(); ++i) pb[i] = cplx(b[i], 0);
     fft(pa, d); fft(pb, d);
     for (int i = 0; i < d; ++i) pa[i] *= pb[i];
     ifft (pa, d);
     vector < int > r(d);
     for (int
           i = 0; i < d; ++i) r[i] = round(pa[i].real());
          while (r.size() \& r.back() = 0) r.pop_back();
     return r;
```

```
Prime
                              Prime
  7681
                   17
                              167772161
                                                  3
  12289
                              104857601
                   11
                                                 3
  40961
                   3
                              985661441
                                                  3
  65537
                   3
                              998244353
                                                  3
                   10
                              1107296257
  786433
  5767169
                   3
                              2013265921
                                                 31
  7340033
                   3
                              2810183681
                                                  11
  23068673
                              2885681153
                   3
                                                 3
  469762049
                   3
                              605028353
                                                 3
7.2 Number Theory Transform* [eleb36]
vector<int> omega;
void Init() {
   omega.resize(kN + 1);
   long long x = fpow(kRoot, (Mod - 1) / kN);
   omega[0] = 1;
for (int i = 1; i <= kN; ++i)
      omega[i] = 1LL * omega[i - 1] * x % kMod;
void Transform(vector<int> &v, int n) {
   \begin{array}{l} BitReverse(v, n);\\ for (int s = 2; s <= n; s <<= 1) \end{array} \{
      int z = s \gg 1;
       for (int i = 0; i < n; i += s) {
          for (int k = 0; k < z; +\!+\!k) {
             \begin{array}{ll} \textbf{int} & x \, = \, 1LL \end{array}
                      * v[i + k + z] * omega[kN / s * k] % kMod;
             v[i + k + z] = (v[i + k] + kMod - x) \% kMod;

(v[i + k] += x) \% = kMod;
  }
void InverseTransform(vector<int> &v, int n) {
    \begin{aligned} & Transform(v, \ n); \\ & for \ (int \ i = 1; \ i < n \ / \ 2; \ +\!\!\!+\!\!\! i) \ swap(v[i], \ v[n \ - i]); \end{aligned} 
   const int kInv = fpow(n, kMod - 2);
   for (int i
           = \ 0; \ i \ < \ n \ ; \ +\!\!\!+\!\! i \ ) \ v \left[ \ i \ \right] \ = \ 1LL \ * \ v \left[ \ i \ \right] \ * \ inv \ \% \ kMod \ ;
7.3 Fast Walsh Transform* [36c9f5]
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) *
invop: or, and, xor = -1, -1, \frac{1}{2} */
void fwt(int *a, int n, int op) { //or
   for (int L = 2; L <= n; L <<= 1)
      for (int i = 0; i < n; i += L)

for (int j = i; j < i + (L >> 1); ++j)

a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f
      N^{!}[1 <\!\!< N] \;,\; g[N][1 <\!\!< N] \;,\; h[N][1 <\!\!< N] \;,\; ct[1 <\!\!< N] \;;
void
       subset_convolution(int *a, int *b, int *c, int L) {
    // c_k = \sum_{i=1}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
   int n = 1 \ll L;
for (int i = 1; i < n; ++i)
      ct[i] = ct[i \& (i - 1)] + 1;
   for (int i = 0; i < n; ++i)
   f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
for (int i = 0; i <= L; ++i)
fwt(f[i], n, 1), fwt(g[i], n, 1);
for (int i = 0; i <= L; ++i)
       for (int j = 0; j \le i; ++j)
   for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];
for (int i = 0; i <= L; ++i)
   fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
      c[i] = h[ct[i]][i];
7.4 Polynomial Operation [37b8c7]
#define
\begin{array}{lll} & \text{fi}\,(s,\,n) \;\; \text{for} \;\; (\text{int}\;i=(\text{int})(s);\;i<(\text{int})(n);\; +\!\!\!+\!\!\!i)\\ & \text{template}{<}\text{int}\;\; \text{MAXN},\; ll\;\; P,\;\; ll\;\; RT>\;\!\!//\;\; \text{MAXN}=2^k\\ & \text{struct}\;\; Poly\;:\;\; \text{vector}{<}ll>\; \{\;\;//\;\; \text{coefficients}\;\; \text{in}\;\; [0\;,\;P) \end{array}
```

using vector<ll>::vector;
static NTTMAXN, P, RT> ntt;

int n() const { return (int)size(); } // n() >= 1

$\label{eq:const_poly} \mbox{Poly}\left(\mbox{const}\ \mbox{Poly}\left(\mbox{const}\ \mbox{Poly}\right.\right.\right. \left.\mbox{int m}\right) \ : \ \mbox{vector}<\mbox{ll}>\!\!(m) \ \left\{\mbox{const}\ \mbox{poly}\right.\right.$ $copy_n(p.data(), min(p.n(), m), data());$ Poly& irev() { return reverse(data(), data() + n()), *this; } Poly& isz(int m) { return resize(m), *this; } Poly& iadd(const Poly &rhs) { // n() == rhs.n() fi(0, n()) if (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P; return *this; Poly& imul(ll k) { fi(0, n()) (*this)[i] = (*this)[i] * k % P;return *this; Poly Mul(const Poly &rhs) const { int m = 1; while (m < n() + rhs.n() - 1) m <<= 1; $\begin{array}{l} {\rm Poly} \ X({}^*{\rm this}\ ,\ m)\ ,\ Y({\rm rhs}\ ,\ m)\ ; \\ {\rm ntt}(X.\,{\rm data}\,()\ ,\ m)\ ,\ {\rm ntt}(Y.\,{\rm data}\,()\ ,\ m)\ ; \\ {\rm fi}\,(\,0\ ,\ m)\ X[\,i\,]\ =\ X[\,i\,]\ *\ Y[\,i\,]\ \%\ P; \end{array}$ ntt(X.data(), m, true); return X. isz(n() + rhs.n() - 1);Poly Inv() const { // (*this)[0] != 0, 1e5/95ms if (n() = 1) return $\{ntt.minv((*this)[0])\};$ int m = 1; while (m < n() * 2) m <<= 1; Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m); Poly Y(*this, m); $ntt\,(\,Xi\,.\,data\,(\,)\,\,,\,\,m)\,\,,\,\,\,ntt\,(\,Y.\,data\,(\,)\,\,,\,\,m)\,\,;$ if ((Xi[i] % P) < 0) Xi[i] += P;ntt(Xi.data(), m, true); return Xi.isz(n()); Poly Sqrt() $const \{ // Jacobi((*this)[0], P) = 1, 1e5/235ms$ if (n()= 1) return {QuadraticResidue((*this)[0], P)}; X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);pair < Poly , Poly > DivMod $\begin{array}{l} (const\ Poly\ \&rhs)\ const\ \{\ //\ (rhs.)back()\ !=\ 0\\ if\ (n()< rhs.n())\ return\ \{\{0\},\ ^*this\};\\ const\ int\ m=n()\ -\ rhs.n()\ +\ 1; \end{array}$ Poly X(rhs); X. irev(). isz(m); Poly Y(*this); Y. irev(). isz(m); $\begin{array}{l} \text{Poly } 1(\ \text{this}),\ 1.11c \ () \ .182 \ (m), \\ \text{Poly } Q = Y.\text{Mul}(X.\text{Inv}()) \ .\text{isz} \ (m) \ .\text{irev} \ (); \\ X = \text{rhs} .\text{Mul}(Q),\ Y = \text{*this}; \\ \text{fi} \ (0,\ n()) \ \ \text{if} \ \ ((Y[i] \ -= X[i]) < 0) \ Y[i] \ += P; \\ \text{return} \ \{Q,\ Y.\text{isz} \ (\text{max}(1,\ \text{rhs}.n() \ -\ 1))\}; \end{array}$ Poly Dx() const { Poly ret(n() - 1); fi(0, ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;return ret.isz(max(1, ret.n())); Poly Sx() const { Poly ret(n() + 1); fi(0, n()) ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;return ret; Poly _tmul(int nn, const Poly &rhs) const { Poly Y = Mul(rhs).isz(n() + nn return $Poly(\hat{Y}.data() + \hat{n}() - 1, Y.data() + Y.n());$ $vector < ll > _eval(const)$ vector<ll> &x, const vector<Poly> &up) const { const int m = (int)x.size(); if (!m) return `{ }; vector < Poly > down(m * 2); $// \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};$ // fi(2, m * (2) down [i] = down [i] / (2). DivMod (up [i]) . second; down[1] = Poly(up[1]).irev().isz(n()).Inv().irev()._tmul(m, *this); fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2]);

```
vector < ll > y(m);
     fi(0, m) y[i] = down[m + i][0];
     return y;
  static vector<Poly> _tree1(const vector<ll> &x) {
    const int m = (int)x.size();
     vector<Poly> up(m * 2);
    fi(0, m) up[m + i] = {(x[i] ? P - x[i] : 0), 1};

for (int i = m - 1; i

> 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
    return up;
  }
  vector
       <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
    auto up = _tree1(x); return _eval(x, up);
  static Poly Interpolate (const vector
       <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
     const int m = (int)x.size();
    fi(0, m) down[m + i] = \{z[i]\};
     for (int i = m - int)
    Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
    return Dx().Mul(Inv()).Sx().isz(n());
  Poly Exp() const { // (*this)[0] = 0, 1e5/360ms
    if (n() = 1) return \{1\};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());Poly Y = X.Ln(); Y[0] = P - 1;
     fi(0, n())
          if'((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;
     return X.Mul(Y).isz(n());
   // M := P(P - 1). If k >= M, k := k \% M + M.
  Poly Pow(ll k) const {
    int nz = 0;
     while (nz < n() \& !(*this)[nz]) ++nz;
    if (nz * min(k, (ll)n()) >= n()) return Poly(n()); if (!k) return Poly(Poly \{1\}, n());
     \begin{array}{l} Poly \ X(data() + nz, \ data() + nz + n() - nz * k); \\ const \ ll \ c = ntt.mpow(X[\ 0\ ], \ k \ \% \ (P - 1)); \end{array} 
    return X.Ln().imul
          (k % P).Exp().imul(c).irev().isz(n()).irev();
  }
  static 11
       LinearRecursion(const vector<ll> &a, const vector
       < ll > &coef, ll n) { // a_n = \sum c_j a_(n-j)}
     const int k = (int)a.size();
     assert((int)coef.size() = k + 1);
    \begin{array}{l} Poly \ C(k+1) \,, \ W(Poly \ \{1\}, \ k) \,, \ M=\{0, \ 1\}; \\ fi(1, \ k+1) \ C[k-i] = coef[i] \,? \ P - coef[i] \,: \, 0; \end{array}
    n \not= 2\,,\; M=M.\,Mul(M)\,.\,DivMod(C)\,.\,second\,;
     fi(0, k) ret = (ret + W[i] * a[i]) \% P;
    return ret;
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template  decltype(Poly_t::ntt) Poly_t::ntt = {};
7.5 Value Polynomial [fad6e7]
struct Poly {
  mint base; // f(x) = poly[x - base]
```

```
vector<mint> poly;
Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
mint get_val(const mint &x) {
  if (x >= base \&\& x < base + SZ(poly))
     return poly[x - base];
  mint rt = 0:
  vector < mint > lmul(SZ(poly), 1), rmul(SZ(poly), 1);
  for (int i = 1; i < \hat{SZ}(poly); ++i)

lmul[i] = lmul[i - 1] * (x - (base + i - 1));
  for (int i = SZ(poly) - 2; i >= 0; --i)

rmul[i] = rmul[i + 1] * (x - (base + i + 1));
  for (int i = 0; i < SZ(poly); ++i)
```

```
return rt;
  void raise() { // g(x) = sigma{base:x} f(x)
if (SZ(poly) == 1 && poly[0] == 0)
      \tt return\,;
    mint nw = get_val(base + SZ(poly));
    poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)
      poly[i] += poly[i - 1];
};
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

 $\label{eq:constant} {\rm for}\beta\, {\rm being some constant.}\, {\rm Polynomial}\, P {\rm such that}\, F(P) = 0 {\rm can}\, {\rm befound it}$ eratively. Denote by Q_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1} \!=\! Q_k \!-\! \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

8 Geometry

Basic [b068f0]

bool same

```
(double a, double b) { return abs(a - b) < eps; }
struct P {
    double x,
   P() : x(0), y(0) \{ \}
  P(): x(u), y(u) {}
P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }
double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
double abs() { return byper(x y); }
    double abs() { return hypot(x, y); }
   P unit() { return *this / abs(); }
   P rot(double o) {
       double c = cos(o), s = sin(o);
return P(c * x - s * y, s * x + c * y);
    double angle() { return atan2(y, x); }
struct L {
   // ax + by + c = 0
    double a, b, c, o;
  P pa, pb;

L(): a(0), b(0), c(0), o(0), pa(), pb() {}

L(P pa, P pb): a(pa.y - pb.y), b(pb.x - pa.x

), c(pa ^ pb), o(atan2(-a, b)), pa(pa), pb(pb) {}

return pa + (pb - pa).unit
   P project(P p) { return pa + (pb - pa).unit
    () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs()) * (pb - pa).abs()); }
};
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
    if (\max(p1.x, p2.x) < \min(p3.x, p4.x) \mid 
            \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
    if (\max(p1.y, p2.y) < \min(p3.y, p4.y))
            \max(p3.y, p4.y) < \min(p1.y, p2.y)) \text{ return false};
   return sign((p3 - p1)
(p4 - p1)) * sign
       (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
       (L x, L y) \{ return same(x.a * y.b, x.b * y.a); \}
        (L x, L y) { return P(-x.b * y.c + x.c * y.b, x .a * y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

8.2 KD Tree [36d550]

```
namespace kdt {
```

```
int root, lc[maxn],
                                                                        // availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
      rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[maxn];
                                                                       vector<P> HPI(vector<L> &ls){
point p[maxn];
int build(int l, int r, int dep = 0) {
                                                                          sort(ls.begin(),ls.end(),cmp)
  if (l = r) return -1;
                                                                          vector < L > pls(1, ls[0]);
  function < bool (const point &, const point
                                                                          for (int i=0; i < (int) ls. size(); ++i) if(!
        &> f = [dep](const point &a, const point &b) {
                                                                               same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
     if (dep \& 1) return a.x < b.x;
                                                                          deque<int> dq; dq.push_back(0); dq.push_back(1);
     else return a.y < b.y;
                                                                       #define meow(a,b,c
  };
                                                                            ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
xl[m] = xr[m] = p[m].x;
yl[m] = yr[m] = p[m].y;
                                                                          for(int i=2;i<(int) pls.size();++i){
  meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();</pre>
                                                                            meow(i,dq[0],dq[1])dq.pop\_front();
                                                                            dq.push_back(i);
  lc[m] = build(1, m, dep + 1);
  if (~lc[m]) {
                                                                          meow(dq
    xl[m] = min(xl[m], xl[lc[m]]);
                                                                               . front(), dq.back(), dq[dq.size()-2])dq.pop\_back();
    xr[m] = max(xr[m], xr[lc[m]]);

yl[m] = min(yl[m], yl[lc[m]]);
                                                                          meow(dq.back(),dq[0],dq[1])dq.pop\_front();
                                                                          if (dq.size()<3u)return vector
    yr[m] = max(yr[m], yr[lc[m]]);
                                                                               <P>(); // no solution or solution is not a convex
  }
                                                                          vector<P> rt;
  rc[m] = build(m + 1, r, dep + 1);
                                                                          \begin{array}{ll} \textbf{for} (\, \textbf{int} \ i \! = \! 0; i \! < \! (\, \textbf{int} \,) \, dq. \, size \, (\,); \! + \! + i \,) \, rt \, . \, push\_back \end{array}
  if (~rc[m]) {
                                                                               (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
    xl[m] = min(xl[m], xl[rc[m]]);
xr[m] = max(xr[m], xr[rc[m]]);
                                                                          return rt;
                                                                       }
    yl[m] = min(yl[m], yl[rc[m]]);
                                                                       8.5 Rotating Sweep Line [b9fa8d]
    yr[m] = max(yr[m], yr[rc[m]]);
                                                                       void rotatingSweepLine(vector<pair<int,int>>> &ps){
  return m;
                                                                          int n=int(ps.size());
                                                                          vector < int > id(n), pos(n);
bool bound(const point &q, int o, long long d) {
                                                                          vector < pair < int, int >> line(n*(n-1)/2);
  double ds = sqrt(d + 1.0);
                                                                          int m=-1
   if \ (q.x < xl[o] - ds \ | \ | \ q.x > xr[o] + ds \ | \ | \\
                                                                          for (int i=0; i< n; ++i) for
       q.y <
                                                                               (int j=i+1;j< n;++j)line[++m]=make\_pair(i,j); ++m;
            yl[o] - ds \mid \mid q.y > yr[o] + ds) return false;
                                                                          sort(line.begin(),line.end(),[&](const
  return true;
                                                                                pair<int, int> &a, const pair<int, int> &b)->bool{
                                                                               if (ps
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
  (a.y - b.y) * 111 * (a.y - b.y);
                                                                                    [a.first].first=ps[a.second].first)return 0;
                                                                               if (ps
                                                                                    [b.first].first=ps[b.second].first)return 1;
                                                                               return (double
void dfs (
                                                                                    ) (ps [a.first].second-ps [a.second].second)/(ps
    const point &q, long long &d, int o, int dep = 0) {
                                                                                    [a.first].first-ps[a.second].first) < (double
  if (!bound(q, o, d)) return;
                                                                                    )(ps[b.first].second-ps[b.second].second
)/(ps[b.first].first-ps[b.second].first);
  long long cd = dist(p[o], q);
if (cd != 0) d = min(d, cd);
  if ((dep & 1)
                                                                          for (int i=0; i< n; ++i) id [i]=i;
    && q.x < p[o].x \mid | !(dep & 1) && q.y < p[o].y) { if (~lc[o]) dfs(q, d, lc[o], dep + 1); if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                          sort(id.begin(),id.end(),[&](const
int &a,const int &b){ return ps[a]<ps[b]; });
                                                                          for(int i=0;i<n;++i)pos[id[i]]=i;
  } else {
     if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                          for (int i=0; i \leqslant m++i)
     if (\sim lc [o]) dfs (q, d, lc [o], dep + 1);
                                                                            auto l=line[i];
                                                                            // meow
                                                                            tie (pos[l.first], pos[l.second],
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
                                                                                  id [pos[l.first]], id [pos[l.second]])=make_tuple
                                                                                  (pos[l.second],pos[l.first],l.second,l.first);
  root = build(0, v.size());
                                                                       }
long long nearest (const point &q) {
  long long res = 1e18;
                                                                       8.6 Triangle Center [33473a]
  dfs\left(q,\ res\,,\ root\,\right);
  return res;
                                                                       Point TriangleCircumCenter(Point a, Point b, Point c) {
                                                                          Point res;
                                                                          8.3 Sector Area [ec8913]
                                                                          double ax = (a.x + b.x) / 2;
  calc area of sector which include a, b
                                                                          double ay = (a.y + b.y) /
double SectorArea(Pa, Pb, double r) {
                                                                          double bx = (c.x + b.x) /
  double o = atan2(a.y, a.x) - atan2(b.y, b.x);
                                                                          double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - b.y) / 2
  while (o \le 0) o += 2 * pi;
while (o >= 2 * pi) o -= 2 * pi;
                                                                         double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)) / (\sin(a1) * \cos(a2) - \sin(a2) * \cos(a1));
return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
  o = min(o, 2 * pi - o);
return r * r * o / 2;
                                                                       Point TriangleMassCenter(Point a, Point b, Point c) {
8.4 Half Plane Intersection [0954c1]
                                                                          return (a + b + c) / 3.0;
bool jizz(L l1, L l2, L l3){
 P p=Intersect(12,13);
  Point\ TriangleOrthoCenter(Point\ a,\ Point\ b,\ Point\ c)\ \{
                                                                          return TriangleMassCenter(a, b
                                                                               , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
bool cmp(const L &a, const L &b){
  return same(
```

Point TriangleInnerCenter(Point a, Point b, Point c) {

Point res;

 $a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa)){>} eps):a.o{<} b.o;$

```
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   double la = len(b - c);
                                                                               p.push\_back(a.c + i.rot(o) * a.r);
                                                                               p.push_back(a.c + i.rot(-o) * a.r);
   double lb = len(a - c);
   double lc = len(a - b);
   res.x = (
                                                                             return p;
       la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = ( la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
                                                                          double IntersectArea(Ca, Cb) {
                                                                             if (a.r > b.r) swap(a, b);
                                                                             double d = (a.c - b.c).abs();
                                                                             if (d >= a.r + b.r - eps) return 0;
                                                                             if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
8.7 Polygon Center [728c3a]
                                                                             double p = acos
                                                                                  ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
Point BaryCenter(vector<Point> &p, int n) {
                                                                             double q = a\cos
  Point res(0, 0);
                                                                             ((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
   double s = 0.0, t;
for (int i = 1; i < p.size() - 1; i++) {
     t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
                                                                          // remove second
     s += t:
                                                                           level if to get points for line (defalut: segment) vector<P> CircleCrossLine(P a, P b, P o, double r) {
     double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y); double C = sq(a.x - o.x)
  res.x /= (3 * s);
res.y /= (3 * s);
                                                                                  ) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  return res;
                                                                             vector<P> t;
                                                                             if (d \ge -eps) {
                                                                               d = \max(0., d);
8.8 Maximum Triangle [55b8cb]
                                                                                double i = (-B - sqrt(d)) / (2 * A);
                                                                                double j = (-B + \operatorname{sqrt}(d)) / (2 * A);
double ConvexHullMaxTriangleArea
                                                                                if (i - 1.0 <= eps && i >=
     (Point p[], int res[], int chnum) {
                                                                                     -eps) t.emplace\_back(a.x + i * x, a.y + i * y);
   double area = 0, tmp;
                                                                                if (j - 1.0 \le eps \& j > =
   res[chnum] = res[0];
                                                                                     -eps) t.emplace_back(a.x + j * x, a.y + j * y);
   for (int i = 0, j = 1, k = 2; i < chnum; i++) {
     while (fabs(Cross(p[
           \begin{array}{l} \text{res}\left[j\right]\right] \text{ - p}\left[\text{res}\left[i\right]\right], \text{ p}\left[\text{res}\left[(k+1)\ \%\ \text{chnum}\right]\right] \text{ - p}\left[\text{res}\left[i\right]\right])) > \text{fabs}\left(\text{Cross}\left(\text{p}\left[\text{res}\left[j\right]\right]\right] \text{ - p}\left[\text{res}\left[i\right]\right]\right), \end{aligned} 
                                                                             return t;
                                                                          // calc area
           p[res[k]] - p[res[i]]))) k = (k + 1) % chnum;
                                                                                 intersect by circle with radius r and triangle OAB
     tmp = fabs(Cross(
                                                                           double AreaOfCircleTriangle(Pa, Pb, double r) {
          p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                                                             bool ina = a.abs() < r, inb = b.abs() < r;
     if (tmp > area) area = tmp;
                                                                             auto p = CircleCrossLine(a, b, P(0, 0), r);
     while (fabs(Cross(p[
           \begin{array}{l} {\rm res}\,[(\,j+1)\,\%\,\,{\rm chnum}\,] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,\, , \,\, p[\,{\rm res}\,[\,k\,]] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,) \,\, > \,\, {\rm fabs}\,(\,{\rm Cross}\,(p[\,{\rm res}\,[\,j\,]] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,\, , \,\, p[\,{\rm res}\,[\,k\,]] \,\, - \,\, p[\,{\rm res}\,[\,i\,]] \,)) \,\, j \,\, = \,\, (j\,+\,1)\,\,\%\,\,{\rm chnum} \,; \\ \end{array} 
                                                                             if (ina) {
                                                                                if (inb) return abs(a ^ b) / 2;
                                                                                return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
     tmp = fabs(Cross(
                                                                             if (inb) return
          p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                                                                   SectorArea(p[0], a, r) + abs(p[0] ^ b) / 2;
     if (tmp > area) area = tmp;
                                                                                 (p.size() = 2u) return SectorArea(a, p[0], + SectorArea(p[1], b, r) + abs(p[0] ^ p[1]
                                                                                                                                    p[1]) / 2;
   return area / 2;
                                                                             else return SectorArea(a, b, r);
8.9 Point in Polygon [88cf80]
                                                                           // for any triangle
                                                                          double AreaOfCircleTriangle(vector<P> ps, double r) {
int pip(vector<P> ps, P p) {
                                                                             double ans = 0;
   int c = 0;
                                                                             for (int i = 0; i < 3; ++i) {
  int j = (i + 1) % 3;
   for (int i = 0; i < ps.size(); ++i) {
     int a = i, b = (i + 1) \% ps.size();
                                                                                double o = atan2
     L l(ps[a], ps[b]);
                                                                               \begin{array}{lll} (ps[\,i\,].\,y,\ ps[\,i\,].\,x) \ \ -\ atan2(\,ps[\,j\,].\,y,\ ps[\,j\,].\,x)\,;\\ if\ (o>=\ pi)\ o=o\ -\ 2\ *\ pi\,;\\ if\ (o<=\ -pi)\ o=o\ +\ 2\ *\ pi\,; \end{array}
     P q = l.project(p);
     if ((p - q).abs() < eps && l.inside(q)) return 1;
     if (same(ps[
                                                                               ans += AreaOfCircleTriangle
          a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
                                                                                     (ps[i], ps[j], r) * (o >= 0 ? 1 : -1);
     if (ps[a].y > ps[b].y) swap(a, b);
     if (ps[a].y <= p.y && p.y <
                                                                             return abs(ans);
            ps[b].y \& p.x \le ps[a].x + (ps[b].x - ps[a].x
          ) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) ++c;
                                                                                    Tangent of Circles and Points to Circle
                                                                          8.11
   return (c & 1) * 2;
                                                                                    [477789]
                                                                           vector<L> tangent(C a, C b) {
                                                                          #define Pij \
8.10 Circle [b6844a]
                                                                             P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
struct C {
                                                                             z.emplace\_back(a.c + i, a.c + i + j);
                                                                          #define deo(I,J)
  Р с;
                                                                             double d = (a
   double r;
  C(P \ c = P(0, 0), \ double \ r = 0) : c(c), r(r) \{\}
                                                                                  .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
                                                                             vector<P> Intersect(C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
                                                                             if (a.r < b.r) swap(a, b);
   vector<P> p;
                                                                             vector<L> z;
   if (same(a.r + b.r,
                                                                             if ((a.c - b.c).abs() + b.r < a.r) return z;
        d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
                                                                             else if (same((a.c - b.c).abs() + b.r, a.r)) \{ Pij; \}
   else if (a.r + b.r > d \&\& d + a.r >= b.r) {
     double o = acos
                                                                               deo(-,+);
```

((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));

P i = (b.c - a.c).unit();

if (same(d, a.r + b.r)) { Pij; }

else if $(d > a.r + b.r) \{ deo(+,-); \}$

8.12 Area of Union of Circles [0590f1]

```
\begin{array}{lll} vector < pair < double \;, & double >>> CoverSegment(C \;\&a, \; C \;\&b) \; \; \{ \\ double \; d = (a.c \; - \; b.c).abs() \;; \end{array}
   \begin{array}{lll} vector < pair < double \,, & double > > res \,; \\ if & (same(a.r + b.r, d)) \;; \end{array} 
  if (same(a.r + b.r, u_{jj}),
else if (d \le abs(a.r - b.r) + eps) {
     if (a.r < b.r) res.emplace_back(0, ...)
  else\ if\ (d < abs(a.r + b.r) - eps) {
     double 1 = z - o, r = z + o;
if (1 < 0) 1 += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
     else res.emplace_back(l, r);
  return res;
double CircleUnionArea
     (vector < C > c)  { // circle should be identical
  int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
  vector<pair<double, double>>> s = {{2 * pi, 9}}, z;
  for (int j = 0; j < n; ++j) if (i != j) {
    z = CoverSegment(c[i], c[j]);
    for (vector for i = 0) a push back(a);
}
        for (auto &e : z) s.push_back(e);
     for (auto &e : s) {
        if (e.first > w) a += F(e.first) - F(w);
        w = \max(w, e.second);
     }
  return a * 0.5;
```

8.13 Minimun Distance of 2 Polygons [e9c988]

8.14 2D Convex Hull [d97646]

```
bool operator < (const P & a, const P & b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator>(const P &a, const P &b) {
  return same(a.x, b.x) ? a.y > b.y : a.x > b.x;
#define crx(a, b, c) ((b - a) (c - a)
vector < P > convex(vector < P > ps)  {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (Pa, Pb) { return
  same(a.x, b.x) ? a.y < b.y : a.x < b.x; }); for (int i = 0; i < ps. size(); ++i) { while (p. size() >= 2 && crx(p[p. size() -
         2], ps[i], p[p.size() - 1]) >= 0) <math>p.pop\_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() -
2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double
      x) \{ return same(x, 0) ? 0 : x > 0 ? 1 : -1; \}
P isLL(P p1, P p2, P q1, P q2) {
  struct CH {
  int n;
  \text{vector} <\!\!P\!\!> p\,,\ u\,,\ d\,;
  CH() {}
  CH(vector < P > ps) : p(ps) {
    n = ps.size();
    \mathtt{rotate}\,(\,\mathtt{p.begin}\,
         ()\;,\;\min\_{element(p.begin(),\;p.end())}\;,\;p.end())\;;
    auto t = max\_element(p.begin(), p.end());
    d = \text{vector} < P > (p. \text{begin}(), \text{next}(t));
    u = \text{vector} < P > (t, p.end()); u.push_back(p[0]);
  int find (vector<P> &v, P d) {
    int l = 0, r = v.size();
    while (1 + 5 < r) {

int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;

if (v[L] * d > v[R] * d) r = R;
       else l = L;
     int x = 1;
     for (int i = l +
           1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
    return x:
  int findFarest(P v) {
     if (v.y > 0 \mid | v.y = 0 & v.x > 0) return
          ((int)d.size() - 1 + find(u, v)) \% p.size();
    return find(d, v);
    get(int 1, int r, Pa, Pb) {
    int s = sgn(crx(a, b, p[1 \% n]));
    while (l + l < r) {
int m = (l + r) >> 1;
        \mbox{if } (sgn(crx(a,\ b,\ p[m\ \%\ n])) == s) \ l = m; \\
    return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
  vector<P> getLineIntersect(P a, P b) {
    int X = findFarest((b - a).rot(pi / 2));
int Y = findFarest((a - b).rot(pi / 2));
    if (X > Y) swap(X, Y);
     if (sgn
    void update_tangent(P q, int i, int &a, int &b) {
```

```
\begin{array}{l} if \ (sgn(crx(q,\ p[a]\,,\ p[i])) > 0) \ a = i\,; \\ if \ (sgn(crx(q,\ p[b],\ p[i])) < 0) \ b = i\,; \end{array}
  void bs(int 1, int r, P q, int &a, int &b) {
     if (l == r) return;
     update_tangent(q, 1 % n, a, b);
     while (l + 1 < r) {
        if (sgn(crx
              else r = m;
     update_tangent(q, r % n, a, b);
  int x = 1;
  for (int i = 1)
         + \ 1; \ i < r\,; \ +\!\!\!+\!\! i\,) \ \ \text{if} \ \ (v\,[\,i\,] \ \ ^* \ d > v\,[\,x\,] \ \ ^* \ d) \ \ x = i\,;
int findFarest(P v) {
  if (v.y > 0 | | v.y = 0 & v.x > 0) return
          ((int)d.size() - 1 + find(u, v)) \% p.size();
  return find(d, v);
Ṕget(int l, int r, Pa, Pb) {
  int s = sgn(crx(a, b, p[1 \% n]));
  while (l + 1 < r) {
     int m = (1 + r) >> 1;
      \mbox{if } (sgn(crx(a,\ b,\ p[m\ \%\ n])) == s) \ l = m; \\
     else r = m;
  return isLL(a, b, p[1 \% n], p[(1 + 1) \% n]);
vector < P > getIS(P a, P b) {
   \begin{array}{ll} \text{int } X = \operatorname{findFarest}((b-a).\operatorname{spin}(pi\ /\ 2));\\ \text{int } Y = \operatorname{findFarest}((a-b).\operatorname{spin}(pi\ /\ 2)); \end{array} 
  if (X > Y) swap(X, Y);
  \begin{array}{c} \text{if } (\operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[X])) \ * \ \operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[Y])) < \\ 0) \ \text{return } \{\operatorname{get}(X,\ Y,\ a,\ b),\ \operatorname{get}(Y,\ X+n,\ a,\ b)\}; \end{array}
  return { };
void update_tangent(P q, int i, int &a, int &b) {
  \begin{array}{ll} \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{a}],\mbox{ }p[\mbox{i}])) > 0) \mbox{ }a = \mbox{i}\,;\\ \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{b}],\mbox{ }p[\mbox{i}])) < 0) \mbox{ }b = \mbox{i}\,; \end{array}
void bs(int 1, int r, Pq, int &a, int &b) {
  if (l == r) return;
  update_tangent(q, 1 % n, a, b);
  while (1 + 1 < r) {
     int m = (l + r) \gg 1;
     if (sgn
           (crx(q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
     else r = m;
  update_tangent(q, r % n, a, b);
bool contain (P p) {
  if (p.x < d[0].x | | p.x > d.back().x) return 0;
  auto it
        = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
  if (it->x == p.x) {
   if (it->y > p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
  i\,t \,=\, lower\_bound
        (u.begin(), u.end(), P(p.x, 1e12), greater < P > ());
      (it -> x == p.x) \{
  if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
  a = b = 0;
  int i
        = lower\_bound(d.begin(), d.end(), p) - d.begin();
  bs(0, i, p, a, b);
  bs(i, d.size(), p, a, b);
i = lower_bound(
        u.begin(), u.end(), p, greater < P > ()) - u.begin();
  bs((int
        d.size() - 1, (int)d.size() - 1 + i, p, a, b);
  bs((int)d.size()
         -1 + i, (int)d.size() - 1 + u.size(), p, a, b);
  return 1;
```

8.15 3D Convex Hull [clae8f]

```
double
      absvol(const P a, const P b, const P c, const P d) {
  struct convex3D {
  static const int maxn=1010;
  struct T{
     int a,b,c;
     bool res;
     T()\{\}
     T(int a, int
           b, int c, bool res=1: a(a), b(b), c(c), res(res)
  int n,m;
  P p[maxn];
  T f [maxn* 8];
  int id [maxn] [maxn];
  bool on (T &t ,P &q) {
     return ((
          p[t.c]-p[t.b])^(p[t.a]-p[t.b])*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
     int g=id[a][b];
     if (f[g].res){
       if (on(f[g],p[q]))dfs(q,g);
       else {
          id [q][b]=id[a][q]=id[b][a]=m;
          f[m++]=T(b,a,q,1);
     }
  }
  void dfs(int p,int i){
     f[i].res=0;
     meow(p, f[i].b, f[i].a);
     meow(p, f[i].c, f[i].b);
     meow(p, f[i].a, f[i].c);
  void operator()(){
     if (n<4)return;
     if([&](){
          \quad \  \  for\,(\,int\ i\!=\!1; i\!<\!\!n;\!+\!\!+\!i\,)\,i\,f\,(\,abs
               (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
          return 1;
}() || [&](){
          for (int i=2; i < n; ++i) if (abs((p[0]-p[i])
                ^(p[1]-p[i]))>eps)return swap(p[2],p[i]),0;
          return 1;
}() || [&](){
          for (int
               =3;i < n;++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))
               *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
          return 1:
          }())return;
     for (int i=0; i<4;++i) {
       T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
       if(on(t,p[i]))swap(t.b,t.c);

id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
       f[m++]=t;
     for (int i=4; i< n; ++i) for
          (int j=0; j < m++j) if (f[j].res & on(f[j],p[i])){
       dfs\left( \,i\,\,,\,j\,\right) ;
       break;
     int mm=m; m=0;
     for (int i=0;i<mm,++i) if (f[i].res) f[m++]=f[i];
  bool same(int i, int j){
     return !(absvol(p[f[i].a],p[f[i
           \begin{array}{l} [\ ] . \ b] . \ p[f[i].c], p[f[j].a]) > eps \ | \ | \ absvol(p[f[i].a], p[f[i].b]) > eps \ | \ | \ absvol(p[f[i].a], p[f[i].b]) > eps \ | \ | \ absvol(p[f[i].a], p[f[i].b], p[f[i].c], p[f[j].c]) > eps); \end{array} 
  int faces(){
     int r=0;
     for (int i=0; i < m,++i) {
       int iden=1;
       for (int j=0; j< i; ++j) if (same(i,j)) iden=0;
       r += iden:
     return r;
```

| } tb;

```
8.16 Minimum Enclosing Circle [7e5b31]
```

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 \hat{p1};
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {
     if (norm2(cent - p[i]) <= r) continue;</pre>
     cent = p[i];
     r = 0.0;
     for (int j = 0; j < i; ++j) {
  if (norm2(cent - p[j]) <= r) continue;
  cent = (p[i] + p[j]) / 2;
        r = norm2(p[j] -
                               cent);
        for (int k = 0; k < j; ++k) {
    if (norm2(cent - p[k]) <= r) continue;
    cent = center(p[i], p[j], p[k]);
           r = norm2(p[k] - cent);
     }
  return circle(cent, sqrt(r));
```

8.17 Closest Pair [7f292a]

```
double closest_pair(int l, int r) {
  // p should be sorted
          increasingly according to the x-coordinates.
  if (l = r) return 1e9;
  if (r - l = 1) return dist(p[l], p[r]);
  int m = (1 + r) >> 1;
  double d =
         min(closest\_pair(l, m), closest\_pair(m + 1, r));
  vector<int> vec;
  for (int i = m; i >= 1 & &
  fabs(p[m].x - p[i].x) < d; --i) vec.push_back(i);
for (int i = m + 1; i <= r &&
    fabs(p[m].x - p[i].x) < d; ++i) vec.push_back(i);
sort(vec.begin(), vec.end()
          [\&](int a, int b) { return p[a].y < p[b].y; });
  for (int i = 0; i < vec.size(); ++i) {
  for (int j = i + 1; j < vec.size()
    && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {
        d = min(d, dist(p[vec[i]], p[vec[j]]));
     }
  return d;
```

9 Else

9.1 Cyclic Ternary Search* [28a883]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n = 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

9.2 Mo's Algorithm(With modification) [5dec12]

```
/*
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
*/
struct Query {
int L, R, LBid, RBid, T;
```

```
 \begin{array}{c} Query(int\ l,\ int\ r,\ int\ t)\colon\\ L(l),\ R(r),\ LBid(l\ /\ blk),\ RBid(r\ /\ blk),\ T(t)\ \{\}\\ bool\ operator<(const\ Query\ \&q)\ const\ \{\\ if\ (LBid\ !=\ q.LBid)\ return\ LBid<\ q.LBid;\\ if\ (RBid\ !=\ q.RBid)\ return\ RBid<\ q.RBid;\\ return\ T<\ b.T;\\ \}\\ \};\\ void\ solve(vector<Query>\ query)\ \{\\ sort(ALL(query));\\ int\ L=0,\ R=0,\ T=-1;\\ for\ (auto\ q:\ query)\ \{\\ while\ (T<\ q.T)\ addTime(L,\ R,\ ++T);\ //\ TODO\\ while\ (T>\ q.T)\ subTime(L,\ R,\ T--);\ //\ TODO\\ while\ (R<\ q.R)\ add(arr[++R]);\ //\ TODO\\ while\ (L>\ q.L)\ add(arr[-L]);\ //\ TODO\\ while\ (R>\ q.R)\ sub(arr[R--]);\ //\ TODO\\ while\ (L<\ q.L)\ sub(arr[L++]);\ //\ TODO\\ while\ (L<\ q.L)\ sub(arr[L++]);\ //\ TODO\\ //\ answer\ query\\ \}\\ \end{array} \right\}
```

9.3 Mo's Algorithm On Tree [4a7f74]

```
Mo's Algorithm On Tree
Preprocess:\\
1) LCA
2)
   dfs with in[u] = dft++, out[u] = dft++
   \operatorname{ord}[\operatorname{in}[u]] = \operatorname{ord}[\operatorname{out}[u]] = u
   bitset ⟨MAXIV inset
4)
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
     int c = LCA(u, v);
     if (c = u \mid \mid c = v)
        q.\,lc\,a\,=\,-1\,,\,\,q.\,L\,=\,out\,[\,c\,\,\,\widehat{}\,\,u\,\,\,\widehat{}\,\,v\,]\,\,,\,\,q.\,R\,=\,out\,[\,c\,\,]\,;
     else if (out[u] < in[v]
        q.lca = c, q.L = out[u], q.R = in[v];
     else
        q.lca = c, q.L = out[v], q.R = in[u];
     q.Lid = q.L / blk;
   bool operator < (const Query &q) const {
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;
void flip(int x) {
     if (inset [x]) sub(arr [x]); // TODO
else add(arr [x]); // TODO
     inset[x] = \sim inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
   int L = 0, R = 0;
   for (auto q : query) {
     while (R < q.R) flip (ord[++R]);
     while (L > q.L) flip (ord[--L]);
     while (R > q.R) flip (ord[R--]); while (L < q.L) flip (ord[L++]);
     if (~q.lca) add(arr[q.lca]);
        answer query
      if (~q.lca) sub(arr[q.lca]);
```

9.4 Additional Mo's Algorithm Trick

- Mo's Algorithm With Addition Only
 - $\ \ Sort querys same as the normal Mo's algorithm.$
 - For each query [l,r]:
 - If l/blk = r/blk, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk + 1) \cdot blk$, curR := curL 1
 - If r > curR, increase curR
 - $\ \operatorname{decrease} \operatorname{\it cur} L \operatorname{tofit} l, \operatorname{and} \operatorname{then} \operatorname{undo} \operatorname{after} \operatorname{answering}$
- $\bullet \quad \text{Mo's Algorithm With Offline Second Time} \\$

 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
 - Part1: Answer all f([1,r],r+1) first.
 - Part2: Store $curR \to R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
 - $-\ \ Note: You must do the above symmetrically for the left boundaries.$

9.5 Hilbert Curve [ed5979]

```
ll hilbert(int n, int x, int y) {
  11 \text{ res} = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx = 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
  return res;
```

9.6 DynamicConvexTrick* [6a6f6d]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
 struct DynamicHull : multiset<Line, less >>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
      lì b) \{ \text{ return a / b - ((a ^ b) < 0 && a % b); } \}
  bool isect(iterator x, iterator y) {
    if (y = end()) \{ x->p = kInf; return 0; \}
    if (x
        ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
    else x->p = Div(y->b - x->b, x->a - y->a);
    return x->p>= y->p;
 f
void addline(11 a, 11 b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
        () && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin
        () && (-x)-p = y-p is ect(x, erase(y));
  il query(ll x) {
   auto l = *lower_bound(x);
    return l.a * x + l.b;
};
```

All LCS* [ae68f0]

```
void all_lcs(string s, string t) { // 0-base
  vector < int > h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {
    int v = -1;
    for (int c = 0; c < SZ(t); ++c)
       if^{\,\,}(s\,[a]\,=\,t\,[\,c\,]\,\,\mid\,|\,\,h\,[\,c\,]\,<\,v)
        swap(h[c], v);
      / LCS(s[0, a], t[b, c]) =
       c - b + 1 - sum([h[i]] >= b] | i <= c)
    // h[i] might become -1 !!
  }
```

AdaptiveSimpson* [dc2085]

```
template<typename Func, typename d = double>
struct Simpson {
  using pdd = pair<d, d>;
  Func f:
 d \ eval(pdd \ l, \ pdd \ r, \ d \ fm, \ d \ eps) \ \{ \ pdd \ m((l.X + r.X) \ / \ 2, \ fm);
    d s = mix(l, r, fm).second;
   auto [flm , sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
    d \ delta = sl + sr - s;
    if (abs(delta
         ) <= 15 * eps) return sl + sr + delta / 15;
    return eval(1, m, flm, eps / 2) +
      eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
```

```
return eval
          (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
   d \text{ eval2}(d l, d r, d \text{ eps}, \text{ int } k = 997) 
     d h = (r - 1) / k, s = 0;

for (int i = 0; i < k; ++i, l += h)

s += eval(l, l + h, eps / k);
     return s;
};
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
      Simulated Annealing [b14262]
double factor = 100000;
const int base = 1e9; // remember to run \sim 10 times
for (int it = 1; it \leq 10000000; ++it) {
     // ans:
          answer, nw: current value, rnd(): mt19937 rnd()
     if (\exp(-(nw - ans))
          ) / factor) >= (double)(rnd() % base) / base)
          ans = nw;
     factor *= 0.99995;
9.10 Tree Hash* [e57357]
ull seed;
ull shift (ull x) {
  x = x << 13;
  x = x >> 7;
  x = x << 17;
   return x;
ull dfs(int u, int f) {
   ull sum = seed;
   for (int i : G[u])
     if (i != f)
       sum += shift(dfs(i, u));
   return sum;
}
9.11 Binary Search On Fraction [951597]
struct Q {
   ll p, q;
  Q go(Q b, ll d) \{ return \{p + b.p*d, q + b.q*d\}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
/// pred(p/q) is true, and 0 \le p,q \le N Q frac_bs(ll N) {
  Q lo\{0, 1\}, hi\{1, 0\};
   if (pred(lo)) return lo;
   assert(pred(hi));
   bool dir = 1, L = 1, H = 1;
for (; L | | H; dir = !dir) {
     11 \text{ len} = 0, \text{ step} = 1;
     for (int t = 0; t < 2 \&\& (t ? step/=2 : step*=2);)
       if (Q mid = hi.go(lo, len + step);
            mid.p > N \mid | \  \, mid.q > N \mid | \  \, \hat{dir} \  \, \hat{\  \, } \  \, pred(mid))
       else len += step;
     swap(lo, hi = hi.go(lo, len));
(dir? L : H) = !!len;
   return dir ? hi : lo;
}
9.12 Bitset LCS [a82d86]
cin >> n >> m;
for (int i = 1, x; i \le n; ++i)
  cin \gg x, p[x].set(i);
for (int i = 1, x; i \le m; i++) {
cin >> x, (g = f) |= p[x];
   f.shiftLeftByOne(), f.set(0);
((f = g - f) ^= g) &= g;
cout << f.count() << '\n';
       Python
10
```

10.1 Misc

from decimal import * setcontext (Context (prec

=MAX_PREC, Emax=MAX_EMAX, rounding=ROUND_FLOOR))

```
print(Decimal(input()) * Decimal(input()))
from fractions import Fraction
Fraction
     ('3.14159').limit_denominator(10).numerator # 22
```