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region and then type : Hash to hash your selection."

"Select

```
1.2 readchar [a419b9]
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
  static char *p = buf, *end = buf;
  if (p == end) end = buf +
       fread_unlocked(buf, 1, bufsize, stdin), p = buf;
1.3 BigIntIO [d9afcb]
      int128 \ x = 0, f = 1;
    char ch = getchar();
    while (ch < '0', | | ch > '9') {
    if (ch == '-') f = -1;
         ch = getchar();
    while (ch >= '0' \& ch <= '9')  {
 x = x * 10 + ch - '0';
         ch = getchar();
void print(__int128 x) {
    if (x < 0) {
putchar('-');
    if (x > 9) print(x / 10);
    putchar(x % 10 + '0'');
bool cmp(\underline{\phantom{a}}int128 x, \underline{\phantom{a}}int128 y) { return x > y; }
1.4 Black Magic [bf0ccd]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
     tree<int, null_type, std::less<int>, rb_tree_tag
      tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int , int> umap;
typedef priority_queue<int> heap;
  mt19937 rng(chrono::
      steady_clock::now().time_since_epoch().count());
      get_rand(int l, int r){ return
        uniform_int_distribution<int>(l, r)(rng); }
  shuffle(v.begin(), v.end(), rng);
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
      (0) = 22; assert(*s.find_by_order(1) = 71);
  assert (s.order_of_key
      (22) = 0; assert (s.order_of_key(71) == 1);
  assert(*s.find_by_order
       (0) = 71); assert(s.order\_of\_key(71) = 0);
  heap a, b; a.join(b);
  std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
```

1.5 Pragma Optimization [eac636]

```
| #pragma GCC optimize("Ofast, no-stack-protector") | #pragma GCC optimize("no-math-errno, unroll-loops") | #pragma GCC target("sse, sse2, sse3, sse3, sse4") | #pragma GCC target("popent, abm, mmx, avx, arch=skylake") | __builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

1.6 Bitset [55712b]

```
#include < bits / stdc + + h>
using namespace std;

int main () {
    bitset <4> bit;
    bit.all(); // all bit is true, ret tru;
    bit.any(); // any bit is true, ret true
    bit.none(); // all bit is false, ret true
    bit.count();
    bit.to_string('0', '1'); // with parmeter
    bit.reset(); // set all to true
    bit.set(); // set all to false
    std::bitset <8> b3{0}, b4{42};
    std::hash<std::bitset <8> hash_fn8(b3); hash_fn8(b4);
}
```

2 Graph

2.1 BCC Vertex* [740acb]

```
struct BCC { // 0-base
   int n, dft, nbcc;
   vector < int > low, dfn, bln, stk, is_ap, cir;
   vector<vector<int>>> G, bcc, nG;
   void make_bcc(int_u) {
      bcc.emplace_back(1, u);
      for (; stk.back() != u; stk.pop_back())
  bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
stk.pop_back(), bln[u] = nbcc++;
   void dfs(int u, int f) {
      int child = 0;
      \begin{array}{ll} low [u] = dfn [u] = ++dft \,, \; stk.pb(u) \,; \\ for \; (int \; v \; : \; G[u]) \end{array}
          if (!dfn[v]) {
             dfs(v, u), ++child;
low[u] = min(low[u], low[v]);
             \begin{array}{l} \mbox{if } (d f n [u] <= low[v]) \ \{ \\ \mbox{is\_ap}[u] = 1, \ b l n [u] = n b c c \,; \end{array}
                make\_bcc(v), bcc.back().pb(u);
         \} \ \ else \ \ if \ \ (dfn[v] < dfn[u] \ \&\& \ v \ != \ f)
      \begin{array}{l} low[u] = min(low[u], dfn[v]); \\ if (f = -1 \&\& child < 2) is_ap[u] = 0; \end{array}
      if (f = -1 \&\& child = 0) make\_bcc(u);
  BCC(int _n): n(_n), dft(),
   G[u].pb(v), G[v].pb(u);
  fvoid solve() {
  for (int i = 0; i < n; ++i)
    if (!dfn[i]) dfs(i, -1);</pre>
   void block_cut_tree() {
      cir.resize(nbcc);
      for (int i = 0; i < n; ++i)
if (is_ap[i])
             bln[i] = nbcc++;
      cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
          for (int j : bcc[i])
             if (is_ap[j])
  \begin{array}{c} n\ddot{G}[\,\bar{i}\,].\,pb(\,bln\,[\,j\,])\,\,,\,\,nG[\,bln\,[\,j\,]].\,pb(\,i\,)\,;\\ \}\,\,//\,\,up\,\,to\,\,2\,\,*\,\,n\,\,-\,\,2\,\,nodes\,!!\,\,\,bln\,[\,i\,]\,\,for\,\,id \end{array}
```

2.2 Bridge* [4da29a]

```
struct ECC { // 0-base
   int n, dft, ecnt, necc;
   vector<int> low, dfn, bln, is_bridge, stk;
   vector<vector<pii>>> G;
   void dfs(int u, int f) {
      dfn[u] = low[u] = ++dft, stk.pb(u);
      for (auto [v, e] : G[u])
```

2.3 SCC* [4057dc]

```
struct SCC { // 0-base
   int n, dft, nscc;
vector<int> low, dfn, bln, instack, stk;
    vector<vector<int>>> G;
    void dfs(int u) {
       low[u] = dfn[u] = ++dft;
       instack[u] = 1, stk.pb(u);
for (int v : G[u])
           if (!dfn[v])
       \begin{array}{l} \operatorname{dfs}(v), \ \operatorname{low}[u] = \min(\operatorname{low}[u], \ \operatorname{low}[v]); \\ \operatorname{else} \ \operatorname{if} \ (\operatorname{instack}[v] \ \&\& \ \operatorname{dfn}[v] < \operatorname{dfn}[u]) \\ \operatorname{low}[u] = \min(\operatorname{low}[u], \ \operatorname{dfn}[v]); \\ \operatorname{if} \ (\operatorname{low}[u] = \operatorname{dfn}[u]) \ \{ \end{array}
           for (; stk.back() != u; stk.pop_back())
              bln[stk
                        .back()] = nscc, instack[stk.back()] = 0;
           instack[u] = 0, bln[u] = nscc++, stk.pop_back();
       }
   SCC(int \underline{n}): n(\underline{n}), dft(), nscc
           (), low(n), dfn(n), bln(n), instack(n), G(n) {}
    void add_edge(int u, int v) {
       G[u].pb(v);
    void solve() {
  for (int i = 0; i < n; ++i)</pre>
           if (!dfn[i]) dfs(i);
}; // scc_id(i): bln[i]
```

2.4 2SAT* [f5630a]

```
struct SAT { // 0-base
   int n;
   vector<bool> istrue;
SCC scc;
SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
   int rv(int a) {
      return a >= n ? a - n : a + n;
   }
   void add_clause(int a, int b) {
      scc.add_edge(rv(a), b), scc.add_edge(rv(b), a);
   }
   bool solve() {
      scc.solve();
      for (int i = 0; i < n; ++i) {
        if (scc.bln[i] == scc.bln[i + n]) return false;
        istrue[i] = scc.bln[i] < scc.bln[i + n];
        istrue[i + n] = !istrue[i];
    }
   return true;
}
</pre>
```

2.5 MinimumMeanCycle* [3e5d2b]

```
ll road[N][N]; // input here
struct MinimumMeanCycle {
    ll dp[N + 5][N], n;
    pll solve() {
        ll a = -1, b = -1, L = n + 1;
    }
}
```

```
for (int i = 2; i <= L; ++i)
        for (int k = 0; k < n; ++k)
           for (int j = 0; j < n; ++j)
dp[i][j] =</pre>
                \min(dp[i - 1][k] + road[k][j], dp[i][j]);
     for (int i = 0; i < n; ++i) {

if (dp[L][i] >= INF) continue;
        11 ta = 0, tb = 1;
for (int j = 1; j < n; +++j)
  if (dp[j][i] < INF &&</pre>
             ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)

ta = dp[L][i] - dp[j][i], tb = L - j;
        if (ta == 0) continue;
        if (a = -1 | | a * tb > ta * b) a = ta, b = tb;
     if (a != -1) {
        11 g = \underline{\hspace{1cm}} gcd(a, b);
        return pll(a / g, b / g);
     return pll(-1LL, -1LL);
   void init(int _n) {
     n\,=\,\underline{\phantom{a}}n\,;
     for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
};
```

2.6 Virtual Tree* [1b641b]

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
   if (top = -1) return st[++top] = u, void();

int p = LCA(st[top], u);

  \begin{array}{l} \mbox{if } (p =\!\!\!\!= st[top]) \mbox{ return } st[++top] = u, \mbox{ void}(); \\ \mbox{while } (top >= 1 \&\& \mbox{ dep}[st[top - 1]] >= \mbox{dep}[p]) \end{array}
  vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
vG[p].pb(st[top]), --top, st[++top] = p;
   st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top =
   \operatorname{sort}\left(\operatorname{ALL}(v)\right)
     [\&](int a, int b) \{ return dfn[a] < dfn[b]; \});
   for (int i : v) insert(i);
   while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
   // do something
   reset(v[0]);
```

Maximum Clique Dyn* [d50aa9]

```
{\tt struct} MaxClique { // fast when N <= 100
 n = _n;
for (int i = 0; i < n; ++i) G[i].reset();
 void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
 void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
   sort (ALL(r)
         , [\&](int x, int y) \{ return d[x] > d[y]; \});
   vector < int > c(SZ(r));
                    q + 1, 1, rgt = 1, tp = 0;
   int lft = max(ans -
   cs[1].reset(), cs[2].reset();
   for (int p : r) {
     int k = 1;
     while ((cs[k] \& G[p]).any()) ++k;
     if (k > rgt) cs[++rgt + 1].reset();
     cs[k][p] = 1;
     if(k < lft) r[tp++] = p;
```

```
\begin{array}{lll} & for \ (int \ k = \ lft \, ; \ k <= \ rgt \, ; \ +\!\!\!+\!\! k) \\ & for \ (int \ p = \ cs \, [\, k\,] \, . \, \_Find\_first \end{array}
           (); p < N; p = cs[k]._Find_next(p))
r[tp] = p, c[tp] = k, ++tp;
      dfs(r, c, l + 1, mask);
   void dfs(vector<</pre>
         int > &r, vector < int > &c, int l, bitset < N > mask) {
      while (!r.empty()) {
         int p = r.back();
         r.pop\_back(), mask[p] = 0;
         if (q + c.back() \le ans) return;
         \operatorname{cur}\left[q++\right] = p;
         vector<int> nr;
         for (int i : r) if (G[p][i]) nr.pb(i);
         if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
         else \ if \ (q > ans) \ ans = q, \ copy\_n(cur \,, \ q, \ sol);
         c.pop_back(), --q;
   int solve() {
      vector < int > r(n);
      ans = q = 0, iota(ALL(r), 0);
      pre_dfs(r, 0, bitset < N > (string(n, '1')));
      return ans;
};
```

2.8 Minimum Steiner Tree* [62d6fb]

```
struct SteinerTree { // 0-base
  \begin{array}{l} int \ n, \ dst\left[N\right]\left[N\right], \ dp\left[1 << T\right]\left[N\right], \ tdst\left[N\right]; \\ int \ vcst\left[N\right]; \ // \ the \ cost \ of \ vertexs \end{array}
   void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i) {
        fill_n (dst[i], n, INF);
        dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
     x = \min(x, val);
   void add_edge(int ui, int vi, int wi) {
     chmin(dst[ui][vi], wi);
   void shortest_path() {
     for (int k = 0; k < n; ++k)
        for (int i = 0; i < n; ++i)
          int solve(const vector<int> ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1; for (int i = 0; i <= full; ++i)
        fill_n (dp[i], n, INF);
     \begin{array}{l} \operatorname{copy\_n(vcst, n, dp[0]);} \\ \operatorname{for (int msk} = 1; \, \operatorname{msk} <= \, \operatorname{full}; \, +\!\!+\!\!\operatorname{msk}) \, \, \{ \end{array}
        if (!(msk & (msk - 1))) {
           int who = 
                          _{
m lg}({
m msk});
           for (int i = 0; i < n; ++i)
             dp [msk
                  [i] = vcst[ter[who]] + dst[ter[who]][i];
        for (int i = 0; i < n; ++i)
           for (int sub = (
                msk - 1) & msk; sub; sub = (sub - 1) & msk)
             chmin (dp [msk] [i],
                  dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
        for (int i = 0; i < n; ++i) {
           tdst[i] = INF;
           for (int j = 0; j < n; ++j)
             chmin(tdst[i], dp[msk][j] + dst[j][i]);
        copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
; // O(V 3^T + V^2 2^T)
```

2.9 Dominator Tree* [2b8b32]

```
struct dominator_tree
                                                                      // 1-base
     vector < int > G[N], rG[N];
     \begin{array}{lll} & \text{int } n, \ pa\left[N\right], \ dfn\left[N\right], \ id\left[N\right], \ Time; \\ & \text{int } semi\left[N\right], \ idom\left[N\right], \ best\left[N\right]; \end{array}
```

```
vector<int> tree[N]; // dominator_tree
   void init (int _n) {
     n = _n;
for (int i = 1; i <= n; ++i)
        G[i].clear(), rG[i].clear();
   void add_edge(int u, int v) {
     G[u].pb(v), rG[v].pb(u);
   void dfs(int u) {
     id[dfn[u] = ++Time] = u;
for (auto v : G[u])
         if \quad (!dfn[v]) \quad dfs(v), \quad pa[dfn[v]] = dfn[u]; 
   int find(int y, int x) {
     if (y <= x) return y;
int tmp = find(pa[y], x);</pre>
     if (semi[best[y]] > semi[best[pa[y]]])
        best[y] = best[pa[y]];
     return pa[y] = tmp;
   void tarjan(int root) {
     Time = 0;
     for (int i = 1; i \le n; ++i) { dfn[i] = idom[i] = 0;
        tree[i].clear();
best[i] = semi[i] = i;
     dfs(root);
     for (int'i = Time; i > 1; --i) {
        int u = id[i];
        for (auto v : rG[u])
          if (v = dfn[v]) {
             tree [semi[i]].pb(i);
        for (auto v : tree[pa[i]]) {
          find(v, pa[i]);
idom[v] =
             semi[best[v]] = pa[i] ? pa[i] : best[v];
        tree [pa[i]]. clear();
     for (int i = 2; i \leftarrow Time; ++i) {
         \begin{array}{ll} \textbf{if} \ (idom\,[\,i\,] \ != \ semi\,[\,i\,]) \ idom\,[\,i\,] \ = \ idom\,[idom\,[\,i\,]]\,; \\ tree\,[\,id\,[idom\,[\,i\,]]\,] \ .pb(\,id\,[\,i\,])\ ; \\ \end{array} 
  }
};
2.10 Four Cycle [584a52]
int main() {
   cin.tie(nullptr)->sync_with_stdio(false);
   cin >> n >> m;
   for (int i = 1; i \le m; i++) {
     int u, v;
     cin >> u >> v;
     E[u].push\_back(v);
     E[v]. push_back(u);
     deg[u]++, deg[v]++;
   for (int u = 1; u \le n; u++)
     for (int v : E[u])

if (\deg[u] > \deg[v] \mid | (\deg[u] \cdot \operatorname{push\_back}(v);

[u] = \deg[v] \&\& u > v) E1[u].push_back(v);
   for (int a = 1; a \le n; a++) {
     for (int b : E1[a])
       for (int c : E[b]) {
  if (deg[a] < deg[c]
                  | | (\deg[a] = \deg[c] \& a \le c)) continue;
           total += cnt[c]++;
     for (int b : E1[a])
        for (int c : E[b]) cnt[c] = 0;
   cout << total << '\n';
   return 0;
}
```

2.11 Minimum Clique Cover* [879472]

```
\begin{array}{lll} & struct & Clique\_Cover \; \{ \; \; // \; \; 0\text{-base} \; , \; \; O(n2^n) \\ & int \; \; co[\; 1 << N] \; , \; \; n \; , \; \; E[N] \; ; \\ & int \; \; dp[\; 1 << N] \; ; \\ & void \; \; init (int \; \_n) \; \; \{ \end{array}
```

```
\begin{array}{l} n = \_n, \ fill\_n \, (dp, \ 1 << n, \ 0) \, ; \\ fill\_n \, (E, \ n, \ 0) \, , \ fill\_n \, (co \, , \ 1 << n, \ 0) \, ; \end{array}
     void add_edge(int u, int v) {
       E[u] = 1 << v, E[v] = 1 << u;
    int solve() {
       for (int i = 0; i < n; +++i)
          co[1 << i] = E[i] | (1 << i);
       co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 -
                                           1:
        for (int i = 1; i < (1 << n); ++i) {
          int t = i & -i;
dp[i] = -dp[i ^ t];
co[i] = co[i ^ t] & co[t];
        for (int i = 0; i < (1 << n); ++i)
           co[i] = (co[i] \& i) == i;
       fwt(co, 1 << n, 1);
for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)</pre>
              sum += (dp[i] *= co[i]);
           if (sum) return ans;
        return n;
};
```

2.12 Number of Maximal Clique * [11fa26]

```
struct BronKerbosch { // 1-base
  \begin{array}{ll} int & n, \ a[N] \ , \ g[N] \ [\grave{N}] \ , \\ int & S, \ all \ [N] \ [N] \ , \ some \ [N] \ [N] \ , \ none \ [N] \ [N] \ ; \end{array}
   void init(int _n) {
     n = 1
      for (int i = 1; i \le n; ++i)
         for (int j = 1; j \le n; ++j) g[i][j] = 0;
  void add_edge(int u, int v) {
     g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
      if (sn = 0 \&\& nn = 0) ++S;
      int u = some[d][0];
      for (int i = 0; i < sn; ++i) {
         int v = some[d][i];
         if (g[u][v]) continue;
        int tsn = 0, tnn = 0
        copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
for (int j = 0; j < sn; ++j)
if (g[v][some[d][j]])
              some[d + 1][tsn++] = some[d][j];
         for (int j = 0; j < nn; +++j)
            if (g[v][none[d][j]])
              none[d + 1][tnn++] = none[d][j];
        \begin{array}{l} dfs\,(d+1,\,an+1,\,tsn\,,\,tnn\,)\,;\\ some\,[d]\,[\,i\,]\,=\,0\,,\,none\,[d]\,[\,nn++]\,=\,v\,; \end{array}
  int solve() {
  iota(some[0], some[0] + n, 1);
     S = 0, dfs(0, 0, n, 0);
      return S;
```

3 Data Structure

3.1 Discrete Trick

3.2 BIT kth* [e39485] | int bit [N + 1]; // $N = 2^k$

```
int query_kth(int k) {
     int res = 0;
     for (int i = N >> 1; i >= 1; i >>= 1)
if (bit[res + i] < k)
              k -= bit[res += i];
     return res + 1;
```

3.3 Interval Container* [c54d29]

```
/* Add and
      remove intervals from a set of disjoint intervals.
  Will merge the added interval with
       any overlapping intervals in the set when adding.
  Intervals are [inclusive, exclusive). */
\operatorname{set} < \operatorname{pii} > ::
     iterator addInterval(set<pii> is, int L, int R) {
  if (L = R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->X \le R) {
 R = \max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y>= L) {
    \begin{array}{l} L = \min(L, \ it \rightarrow X); \\ R = \max(R, \ it \rightarrow Y); \end{array}
    is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii> is, int L, int R) {
 if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it ->Y;
  if (it->X == L) is.erase(it);
  else (int\&)it -> Y = L;
  if (R!= r2) is.emplace(R, r2);
```

3.4 Leftist Tree [e91538]

```
struct node {
  ll v, data; sz, sum; node *1, *r;
   node(ll k)
     v(0), data(k), sz(1), l(0), r(0), sum(k) {}
il sz(node *p) { return p ? p->sz : 0; }
Il sz(node 'p) { return p ! p->sz : 0; }

Il V(node *p) { return p ! p->v : -1; }

Il sum(node *p) { return p ! p->sum : 0; }

node *merge(node *a, node *b) {
  if (!a | | !b) return a ? a : b;
if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
  return a;
void pop(node *&o) {
 node *tmp = o;
  o = merge(o->l, o->r);
  delete tmp:
```

Heavy light Decomposition* [b004ae]

```
ruct Heavy_light_Decomposition { // 1-base int n, ulink [N], deep [N], mxson [N], w[N], pa[N];
struct Heavy_light_Decomposition {
  int t, pl[N], data[N], val[N]; // val: vertex data vector<int> G[N];
  void init(int _n) {
    n = _n;
for (int i = 1; i <= n; ++i)
      G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b) {
    G[a].pb(b), G[b].pb(a);
  void dfs(int u, int f, int d) {
    w[u] = 1, pa[u] = f, deep[u] = d++;
    for (int &i : G[u])
       if (i != f)
         dfs(i, u, d), w[u] += w[i];
         if (w[mxson[u]] < w[i]) mxson[u] = i;
  }
```

```
\begin{array}{ll} void \;\; cut(int\;\; u,\;\; int\;\; link) \;\; \{ \\ \;\; data[\,pl\,[u]\, = +\!\!+\!\!t\,] \;\; = \; val\,[u] \;, \;\; ulink\,[u] \;= \; link \,; \end{array}
      if (!mxson[u]) return;
      cut(mxson[u], link);
for (int i : G[u])
         if (i != pa[u] && i != mxson[u])
  cut(i, i);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
      int ta = ulink[a], tb = ulink[b], res = 0;
      while (ta != tb) {
         if (deep
                [ta] > deep[tb] swap(ta, tb), swap(a, b);
          // query(pl[tb], pl[b])
         tb = ulink[b = pa[tb]];
       if (pl[a] > pl[b]) swap(a, b); 
      // query(pl[a], pl[b])
};
```

3.6 Centroid Decomposition* [5a24da]

```
struct Cent_Dec { // 1-base
  vector < pll > G[N];
  pll info[N]; // store info. of itself
  pll upinfo[N]; // store info. of climbing up
                    layer[N], sz[N], done[N];
   int n, pa[N],
  ll \ dis [\underline{\hspace{1cm}} lg(N) + 1][N];
  void init(int _n) {
  n = _n, layer[0] = -1;
     fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
     for (int i = 1; i \le n; ++i) G[i].clear();
  void add_edge(int a, int b, int w) {
    G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
     int mxsz = 0;
     sz[u] = 1;
     for (pll e : G[u])
       if (!done[e.X] && e.X != f) {
  get_cent(e.X, u, mx, c, num);
          sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
     if(mx > max(mxsz, num - sz[u]))
       mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
   // if required, add self info or climbing info
   dis[layer[org]][u] = d;
   for (pll e : G[u])
       if (!done[e.X] && e.X != f)
          dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
     int mx = 1e9, c = 0, lc;
     \mathtt{get\_cent}\,(\mathtt{u}\,,\ \mathsf{f}\,,\ \mathtt{mx},\ \mathsf{c}\,,\ \mathtt{num})\,;
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
for (pll e : G[c])
       if (!done[e.X]) {
          if (sz[e.X] > sz[c])
          \begin{array}{ll} lc = cut(e.X, c, num - sz[c]); \\ else \ lc = cut(e.X, c, sz[e.X]); \end{array}
          upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
void modify(int u) {
     for (int a = u, ly = layer[a]; a;
                         --ly)
           a = pa[a],
        info[a].X += dis[ly][u], ++info[a].Y;
       if (pa[a])
          upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
  11 query(int u) {
    11 \text{ rt} = 0;
     for (int a = u, ly = layer[a]; a;
       a = pa[a], --ly) {
rt += info[a].X + info[a].Y * dis[ly][u];
       if (pa[a])
```

upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];

rt -=

```
for (Splay *q = x;; q = q->f) {
    return rt;
                                                                          splayVec.pb(q);
  }
                                                                          if (q->isr()) break;
                                                                       reverse (ALL(splayVec));
      LiChaoST* [4a4bee]
                                                                       for (auto it : splayVec) it->push();
                                                                       while (!x->isr()) {
struct L {
                                                                          if (x->f->isr()) rotate(x);
  ll m, k, id;
                                                                          else if (x->dir() = x->f->dir())
  L() : id(-1) \{ \}
                                                                            rotate(x->f), rotate(x);
   \begin{array}{l} L(11\ a,\ 11\ b,\ 11\ c)\ :\ m(a)\ ,\ k(b)\ ,\ id\ (c)\ \{\}\\ 11\ at\ (11\ x)\ \{\ return\ m\ ^*\ x\ +\ k;\ \} \end{array} 
                                                                          else rotate(x), rotate(x);
class LiChao { // maintain max
                                                                     Splay* access(Splay *x) {
private:
                                                                       Splay *q = nil;
 int n; vector<L> nodes;
                                                                       for (; x != nil; x = x->f)
  \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
    if (nodes[rt].id == -1)
      return nodes[rt] = ln, void();
                                                                     void root_path(Splay *x) { access(x), splay(x); }
    bool atLeft = nodes [rt]. at(1) < ln.at(1); if (nodes [rt]. at(m) < ln.at(m))
                                                                     void chroot(Splay *x){
                                                                       root_path(x), x->give_tag(1);
                = 1, swap(nodes[rt], ln);
                                                                       x->push(), x->pull();
    if (r - l = 1) return;
    if (atLeft) insert(l, m, rt << 1, ln);</pre>
                                                                     void split (Splay *x, Splay *y) {
    else insert (m, r, rt << 1 | 1, ln);
                                                                       chroot(x), root_path(y);
  ll query(int l, int r, int rt, ll x) {
                                                                     void link(Splay *x, Splay *y) {
    int m = (1 + r) >> 1; ll ret = -INF;
if (nodes[rt].id != -1) ret = nodes[rt].at(x);
                                                                       root_path(x), chroot(y);
x->setCh(y, 1);
    if (r - l == 1) return ret;
    if (x
                                                                     void cut(Splay *x, Splay *y) {
         < m) return max(ret, query(1, m, rt << 1, x));
                                                                       split(x, y);
    return max(ret, query(m, r, rt \ll 1 | 1, x));
                                                                       if (y->size != 5) return;
                                                                       y->push();
public:
                                                                       y\text{-}\!>\!ch[\,0\,]\,=\,y\text{-}\!>\!ch[\,0\,]\text{-}\!>\!f\,=\,n\,i\,l\,;
 LiChao(int n_) : n(n_), nodes(n * 4) {}
  void insert(L ln) { insert(0, n, 1, ln); }
                                                                     Splay* get_root(Splay *x)
  ll query(ll x) { return query(0, n, 1, x); }
                                                                       for (\text{root\_path}(x); x->\text{ch}[0] != \text{nil}; x = x->\text{ch}[0])
                                                                         x->push();
                                                                       splay(x);
3.8 Link cut tree* [a35b5d]
                                                                       return x;
struct Splay { // xor-sum
                                                                     bool conn(Splay *x, Splay *y) {
  static Splay nil;
                                                                       return get\_root(x) == get\_root(y);
  Splay *ch[2], *f;
  int val, sum, rev, size;
                                                                     Splay* lca(Splay *x, Splay *y) {
  Splay (int
                                                                       access(x), root_path(y);
        _{\text{val}} = 0: val(_{\text{val}}), sum(_{\text{val}}), rev(0), size(1)
                                                                       if (y->f = nil) return y;
  \{ f = ch[0] = ch[1] = &nil; \}
                                                                       return y->f;
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
                                                                     void change(Splay *x, int val) {
  int dir()
                                                                       splay(x), x->val = val, x->pull();
            f->ch[0] = this ? 0 : 1; 
  void setCh(Splay *c, int d) {
                                                                     int query(Splay *x, Splay *y) {
    ch[d] = c;
if (c != &nil) c->f = this;
                                                                       split(x, y);
                                                                       return y->sum;
    pull();
  void give_tag(int r) {
                                                                     3.9 KDTree [375ca2]
    if (r) \overline{\text{swap}}(\text{ch}[0], \text{ch}[1]), rev = 1;
                                                                     namespace kdt {
  void push() {
  if (ch[0] != &nil) ch[0]->give_tag(rev);
                                                                     int root, lc [maxn], rc [maxn], xl [maxn], xr [maxn],
                                                                       yl[maxn], yr[maxn];
                                                                     point p[maxn];
int build(int l, int r, int dep = 0) {
if (l = r) return -1;
    if (ch[1] != &nil) ch[1]->give_tag(rev);
    rev = 0;
  void pull() {
                                                                       function < bool (const point &, const point &)> f =
    // take care of the nil!
                                                                          [dep](const point &a, const point &b) {
    size = ch[0]->size + ch[1]->size + 1;

sum = ch[0]->sum ^ ch[1]->sum ^ val;

if (ch[0]!= &nil) ch[0]->f = this;
                                                                            if (dep \& 1) return a.x < b.x;
                                                                            else return a.y < b.y;
                                                                       int m = (l + r) \gg 1;

nth\_element(p + l, p + m, p + r, f);
    if (ch[1] != &nil) ch[1] -> f = this;
} Splay::nil;
                                                                       xl[m] = xr[m] = p[m].x;
Splay * nil = &Splay:: nil;
void rotate(Splay *x) {
                                                                       yl[m] = yr[m] = p[m].y;
                                                                       lc[m] = build(1, m, dep + 1);
  Splay *p = x - f;
                                                                       if (~lc[m]) {
  int d = x->dir();
                                                                          xl[m] = min(xl[m], xl[lc[m]]);
  if (!p->isr()) p->f->setCh(x, p->dir());
                                                                          xr[m] = max(xr[m], xr[lc[m]]);
                                                                          yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
  else x->f = p->f;
  p-\operatorname{setCh}(x-\operatorname{ch}[!d], d);
  x->setCh(p, !d);
 p->pull(), x->pull();
                                                                       rc[m] = build(m + 1, r, dep + 1);
                                                                       if (~rc[m]) {
                                                                          xl[m] = min(xl[m], xl[rc[m]]);
void splay (Splay *x) {
  vector<Splay*> splayVec;
                                                                          xr[m] = max(xr[m], xr[rc[m]]);
```

```
\begin{array}{l} {\rm yl}\,[m] \; = \; \min(\,{\rm yl}\,[m] \;, \; \; {\rm yl}\,[\,{\rm rc}\,[m]\,]\,) \;; \\ {\rm yr}\,[m] \; = \; \max(\,{\rm yr}\,[m] \;, \; \; {\rm yr}\,[\,{\rm rc}\,[m]\,]\,) \;; \end{array}
                                                                              bool erase(node *&o, int k) {
                                                                                if (!o) return 0;
if (o->data == k) {
  return m;
                                                                                   node *t = o;
bool bound(const point &q, int o, long long d) {
                                                                                   o \rightarrow down(), o = merge(o \rightarrow l, o \rightarrow r);
  double ds = \operatorname{sqrt}(d + 1.0);
                                                                                   delete t;
  return 1;
     return false;
                                                                                 node *\&t = k < o->data ? o->l : o->r;
  return true;
                                                                                 return erase(t, k) ? o->up(), 1 : 0;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
                                                                              void insert(node *&o, int k) {
                                                                                 \quad node \ ^*a\,, \ ^*b\,;
                                                                                 split(o, a, b, k),
                                                                                   o = merge(a, merge(new node(k), b));
void dfs (
  const point &q, long long &d, int o, int dep = 0) {
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x | |
                                                                              void interval(node *&o, int l, int r) {
                                                                                node *a, *b, *c;
split2(o, a, b, l - 1), split2(b, b, c, r);
                                                                                 // operate
                                                                                 o = merge(a, merge(b, c));
     Flow/Matching
  } else {
                                                                              4.1 Dinic [98fb3a]
     if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
     if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                              {\tt struct} \  \, {\tt MaxFlow} \  \, \{ \  \, // \  \, {\tt 0-base}
                                                                                 struct edge {
                                                                                   int to, cap, flow, rev;
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
                                                                                 {\tt vector}{<\!\!\operatorname{edge}\!\!>}\; G[M\!A\!X\!N]\,;
  root = build(0, v.size());
                                                                                 int s, t, dis [MAXN], cur [MAXN], n;
                                                                                 int dfs(int u, int cap) {
long long nearest (const point &q) {
                                                                                   \begin{array}{ll} \mbox{if } (u =\!\!\!\!= t \mid \mid ! \, cap) \ \ \mbox{return cap}; \\ \mbox{for } (\mbox{int } \&i = cur[u]; \ i < (\mbox{int})G[u]. \, size(); \ +\!\!\!\!\!+i) \end{array} \}
  long long res = 1e18;
  dfs(q, res, root);
                                                                                      edge \&e = G[u][i];
  return res;
                                                                                      if (dis[e.to] = dis[u] + 1 & e.flow != e.cap) {
                                                                                         int df = dfs(e.to, min(e.cap - e.flow, cap));
} // namespace kdt
                                                                                         if (df) {
                                                                                           e.flow += df;
3.10 \quad Treap \ [5ab1a1]
                                                                                           G[e.to][e.rev].flow -= df;
struct node {
                                                                                           return df;
  int data, sz;
node *l, *r;
                                                                                      }
  node(int \ k) : data(k), sz(1), l(0), r(0) \{ \}
  void up() {
                                                                                    dis[u] = -1;
                                                                                   return 0;
     if (1) sz += 1->sz;
     if (r) sz += r->sz;
                                                                                 bool bfs() {
                                                                                   fill_n (dis, n, -1);
                                                                                   queue<int> q;
q.push(s), dis[s] = 0;
  void down() {}
f,
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
   if (!a || !b) return a ? a : b;
                                                                                    while (!q.empty()) {
                                                                                      int tmp = q.front();
                                                                                      q.pop();
  if \ (rand() \% (sz(a) + sz(b)) < sz(a))
                                                                                      for (auto &u : G[tmp])
     return a->down(), a->r = merge(a->r, b), a->up(),
                                                                                         if (!~dis[u.to] && u.flow != u.cap) {
                                                                                           q.push(u.to);
                                                                                            dis[u.to] = dis[tmp] + 1;
  return b->down(), b->l = merge(a, b->l), b->up(), b;
void split (node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
                                                                                   return dis[t] != -1;
  o->down();
                                                                                 int maxflow(int _s, int _t) {
  if (o->data \le k)
     a = o, split (o->r, a->r, b, k), a->up();
                                                                                   s = \underline{\ }s, \ t = \underline{\ }t;
int flow = 0, df;
  else b = o, split (o > l, a, b > l, k), b > up();
                                                                                    while (bfs()) {
void split2 (node *o, node *&a, node *&b, int k) {
                                                                                      fill_n(cur, n, 0);
                                                                                      while ((df = dfs(s, INF))) flow += df;
  if (sz(o) \le k) return a = o, b = 0, void();
  o > down();
  if (sz(o->1) + 1 \le k)
                                                                                   return flow;
     a = o, split2(o->r, a->r, b, k - sz(o->l) - 1);
  else b = o, split2(o->l, a, b->l, k);
                                                                                 void init(int _n) {
                                                                                   o > up();
node *kth(node *o, int k) {
  if (k \le sz(o->l)) return kth(o->l, k);
                                                                                 void reset() {
  if (k = sz(o->l)'+1) return o;
                                                                                   for (int i = 0; i < n; ++i)
                                                                                      for (auto \&j : G[i]) j.flow = 0;
  return kth(o->r, k - sz(o->l) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
                                                                                 void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, (int)G[v].size()});
  G[v].pb(edge{u, 0, 0, (int)G[u].size() - 1});
  if (o->data < key)
     return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->l, key);
                                                                             };
```

4.2 Bipartite Matching* [784535]

```
struct Bipartite_Matching { // 0-base
  vector < int > G[N + 1];
  bool dfs(int u) {
    int e = G[u][i];
      if (mq[e] == 1
            || (dis[mq[e]] = dis[u] + 1 \& dfs(mq[e]))|
         return mp[mq[e] = u] = e, 1;
    return dis[u] = -1, 0;
  bool bfs() {
    queue<int> q;
    fill_n (dis, l + 1, -1);
for (int i = 0; i < l; ++i)
      if (!~mp[i])
        q.push(i), dis[i] = 0;
    while (!q.empty())
      int u = q.front();
      q.pop();
      for (int e : G[u])
        if (!~dis[mq[e]])
          q.push(mq[e]), dis[mq[e]] = dis[u] + 1;
    return dis[1] != -1;
  int matching() {
    int res = 0;
    fill\_n\left(mp,\ l\ ,\ -1\right),\ fill\_n\left(mq,\ r\ ,\ l\ \right);
    while (bfs()) {
      fill_n(cur, 1, 0);

for(int i = 0; i < 1; ++i)
        res += (! \sim mp[i] \&\& dfs(i));
    return res; // (i, mp[i] != -1)
  void add_edge(int s, int t) { G[s].pb(t); }
  void init(int _l, int _r) {
    l = \underline{l}, r = \underline{r};
    for (int i = 0; i \le l; ++i)
      G[i].clear();
};
```

4.3 Kuhn Munkres* [4b3863]

```
struct KM { // 0-base, maximum matching
 11 w[N] [N], h1 [N], hr [N], slk [N]; int f1 [N], fr [N], pre [N], qu [N], ql, qr, n; bool v1 [N], vr [N];
  void init(int _n) {
   void add_edge(int a, int b, ll wei) {
   w[a][b] = wei;
  bool Check(int x) {
if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0:
  void bfs(int s) {
    fill_n (slk
        , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] && slk
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
         else slk[x] -= d;
        if (vr[x]) hr[x] = d;
```

4.4 MincostMaxflow* [1c78db]

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
     ll from, to, cap, flow, cost, rev;
   } *past[N];
   vector < Edge > G[N];
  \begin{array}{ll} int & inq\left[N\right], \ n, \ s, \ t; \\ 11 & dis\left[N\right], \ up\left[N\right], \ pot\left[N\right]; \\ bool & BellmanFord\left(\right) \ \{ \end{array}
     fill_n(dis, n, INF), fill_n(inq, n, 0);
     queue<int> q;
     if (!inq[u]) inq[u] = 1, q.push(u);
       }
     };
     relax(s, 0, INF, 0);
while (!q.empty()) {
       int u = q.front();
        q.pop(), inq[u] = 0;
       for (auto &e : G[u]) {
          11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
                (e.to, d2, min(up[u], e.cap - e.flow), &e);
       }
     return dis[t] != INF;
  void solve(int
       , int _t, ll &flow, ll &cost, bool neg = true) {
     s = _s, t = _t, flow = 0, cost = 0;
     if (neg) BellmanFord(), copy_n(dis, n, pot);
     for (; BellmanFord(); copy_n(dis, n, pot)) {
       for (int
       i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
        for (int i = t; past[i]; i = past[i]->from) {
          auto &e = *past[i]
          e.flow += up[t], G[e.to][e.rev].flow -= up[t];
    }
  }
  void init(int _n) {
     n = \underline{n}, fill_n(pot, n, 0);
     for (int i = 0; i < n; ++i) G[i].clear();
  void add_edge(ll a, ll b, ll cap, ll cost) {
   G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
   G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
```

4.5 Maximum Simple Graph Matching* [0fe1c3]

```
struct Matching { // 0-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
        if (vis[x] == tk) return x;
        vis[x] = tk;
        x = Find(pre[match[x]]);
    }
  }
  void Blossom(int x, int y, int l) {
```

```
for (; Find(x) != 1; x = pre[y]) {
    pre[x] = y, y = match[x];
if (s[y] = 1) q.push(y), s[y] = 0;
    for (int z: \{x, y\}) if (fa[z] = z) fa[z] = 1;
 }
bool Bfs(int r) {
  iota(ALL(fa), 0); fill(ALL(s), -1);
  q = queue<int>(); q.push(r); s[r] = 0;
for (; !q.empty(); q.pop()) {
    for (int x = q. front(); int u : G[x])
      if(s[u] = -1) {
         last =
                 match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]); s[match[u]] = 0;
      } else if (!s[u] \&\& Find(u) != Find(x)) {

int l = LCA(u, x);

         Blossom(x, u, l); Blossom(u, x, l);
  return false;
Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis
(n+1), \text{ pre}(n+1, n), \text{ match}(n+1, n), G(n) \{\}
\text{void add\_edge}(\text{int } u, \text{ int } v)
{ G[u].pb(v), G[v].pb(u); }
int solve() {
  int ans = 0;
  for (int x = 0; x < n; ++x)
     if (match[x] == n) ans += Bfs(x); 
  return ans:
} // match[x] == n means not matched
```

4.6 Maximum Weight Matching* [1ec446]

```
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
      struct edge { int u, v, w; }; int n, nx;
       vector<int> lab; vector<vector<edge>>> g;
       \begin{array}{l} \text{vector} < \text{int} > \text{slk} \;,\;\; \text{match} \;,\; \text{st} \;,\; \text{pa} \;,\; S \;,\; \text{vis} \;; \\ \text{vector} < \text{vector} < \text{int} > \; \text{flo} \;,\;\; \text{flo\_from} \;;\;\; \text{queue} < \text{int} > \; q \;; \\ \text{WeightGraph} (\;\; \text{int} \;\; \text{n}\_) \;:\;\; \text{n}(\text{n}\_) \;,\;\; \text{nx}(\text{n} \;\;^* \;\; 2) \;,\;\; \text{lab}(\text{nx} \;+\; 1) \;,\;\; \\ \text{g}(\text{nx} \;+\; 1 \;,\;\; \text{vector} < \text{edge} > (\text{nx} \;+\; 1) \;,\;\; \text{slk} (\text{nx} \;+\; 1) \;,\;\; \\ \end{array} 
             flo (nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
match = st = pa = S = vis = slk;
            REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
       int E(edge e)
      { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; } void update_slk(int u, int x, int &s) { if (!s || E(g[u][x]) < E(g[s][x])) s = u; }
       void set_slk(int x) {
               slk[x] = 0;
                      \begin{tabular}{ll} \be
                            update\_slk(u, x, slk[x]);
       void q_push(int x)
              if (x \le n) q.push(x);
              else for (int y : flo[x]) q_push(y);
       void set_st(int x, int b) {
              st[x] = b;
              if (x > n) for (int y : flo[x]) set_st(y, b);
       vector<int> split_flo(auto &f, int xr) {
              auto it = find(ALL(f), xr);
              if (auto pr = it - f.begin(); pr % 2 == 1)
                     reverse(1 + ALL(f)), it = f.end() - pr;
              auto res = vector(f.begin(), it);
              return f.erase(f.begin(), it), res;
       void set_match(int u, int v) {
            match[u] = g[u][v].v;
              if (u \le n) return;
              \begin{array}{ll} \textbf{int} & \textbf{xr} = \textbf{flo\_from} \, [\, \textbf{u} \,] \, [\, \textbf{g} \, [\, \textbf{u} \,] \, [\, \textbf{v} \,] \, . \, \textbf{u} \,] \,; \end{array}
            set_match(xr, v); f.insert(f.end(), ALL(z));
       void augment(int u, int v) {
```

```
for (;;) {
    int xnv = st[match[u]]; set\_match(u, v);
    if (!xnv) return;
    set_match(v = xnv, u = st[pa[xnv]]);
int lca(int u, int v) {
  static int t = 0; ++t;
  for (++t; u \mid | v; swap(u, v)) if (u) {
    if (vis[u] == t) return u;
    vis[u] = t, u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0, match[b] = match[o];
  vector < int > f = \{o\};
  for (int t : {u, v}) {
    reverse(1 + ALL(f));
    f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set\_st(b, b);
 REP(x, 1, nx) g[b][x].w = g[x][b].w = 0; fill (ALL(flo_from[b]), 0);
  for (int xs : flo[b]) {
    REP(x, 1, nx)
       if (g[b][x].w = 0 \mid | E(g[xs][x]) < E(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
            1, n
       if (flo_from [xs][x]) flo_from [b][x] = xs;
  set_slk(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
  for (int x: split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
    slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
  for (int x : flo[b])

if (x = xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set\_slk(x);
  st[b] = 0;
bool on_found_edge(const edge &e) {
   \begin{array}{lll} & \text{if (int } u = st \, [e \cdot u] \;, \; v = st \, [e \cdot v] \;; \; S[v] = -1) \; \{ \\ & \text{int } nu = st \, [match[v]] \;; \; pa[v] = e \cdot u \;; \; S[v] = 1 \;; \\ & \text{slk} \, [v] = slk \, [nu] = S[nu] = 0 \;; \; q\_push(nu) \;; \\ \end{array} 
  else if (S[v] = 0)
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
  fill(ALL(S), -1), fill(ALL(slk), 0);
    = queue < int > ();
  for (;;)
    while (SZ(q)) {
      int u = q.front(); q.pop();
       if (S[st[u]] = 1) continue;
      REP(\,v\,,\ 1\,,\ n\,)
         update\_slk(u, st[v], slk[st[v]]);
           else if
                 (on_found_edge(g[u][v])) return true;
    int d = INF;
    REP(x, 1, nx)
       if (int
            s = slk[x]; st[x] == x \&\& s \&\& S[x] <= 0
         d = min(d, E(g[s][x]) / (S[x] + 2));
    REP(u, 1, n)
       if (S[st[u]] == 1) lab[u] += d;
```

void init(int _n) {

 $n = _n;$ for (int i = 0; i < n + 2; ++i)

G[i].clear(), cnt[i] = 0;

```
\begin{array}{ll} \text{else if } (S[st[u]] =\!\!\!\!= 0) \ \{ \\ \text{if } (lab[u] <\!\!\!= d) \ \text{return false}; \end{array}
                                                                                    \begin{array}{l} void \ add\_edge(int \ u, \ int \ v, \ int \ lcap \,, \ int \ rcap) \ \{ \\ cnt[u] \ -= \ lcap \,, \ cnt[v] \ += \ lcap \,; \end{array} 
                                                                                      G[\,u\,] \cdot pb\big(\,edge\big\{v\,,\ rcap\,,\ lcap\,,\ SZ(G[\,v\,]\,)\,\big\})\,;
             lab[u] -= d;
                                                                                      G[v].pb(edge\{u, 0, 0, SZ(G[u]) -
        void add_edge(int u, int v, int cap) { G[u].pb(edge\{v, cap, 0, SZ(G[v])\});
        REP(x, 1, nx)
           if (int s = slk[x]; st[x] == x &&
                                                                                      G[v].pb(edge\{u, 0, 0, SZ(G[u]) - 1\});
                s \&\& st[s] != x \&\& E(g[s][x]) == 0
              if (on_found_edge(g[s][x])) return true;
                                                                                   int dfs(int u, int cap) {
                                                                                       \begin{array}{lll} & \text{if } (u = t \mid \mid ! \operatorname{cap}) \text{ return } \operatorname{cap}; \\ & \text{for } (\operatorname{int \&i} = \operatorname{cur}[u]; \text{ } i < \operatorname{SZ}(G[u]); \text{ } +\!\!+\!\! i) \end{array} \{ 
        REP(b, n + 1, nx)

if (st[b] == b && S[b] == 1 && lab[b] == 0)
                                                                                         edge &e = G[u][i];
if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
              expand_blossom(b);
                                                                                            int df = dfs(e.to, min(e.cap - e.flow, cap));
     return false;
                                                                                            if (df) {
   }
   pair<ll, int> solve() {
                                                                                              e.flow += df, G[e.to][e.rev].flow -= df;
     fill(ALL(match), 0);
                                                                                              return df;
     REP(u, 0, n) st [u] = u, flo [u]. clear();
                                                                                        }
     int w_max = 0;
     dis[u] = -1;
        w_{max} = max(w_{max}, g[u][v].w);
                                                                                      return 0;
     fill (ALL(lab), w_max);
int n_matches = 0; 11 tot_weight = 0;
                                                                                   bool bfs() {
                                                                                      fill_n(dis, n + 3, -1);
      while (matching()) ++n_matches;
                                                                                      queue<int> q;
     REP(u, 1, n) \ if \ (match[u] \ \&\& \ match[u] < u)
                                                                                      q.push(s), dis[s] = 0;
        tot_weight += g[u][match[u]].w;
                                                                                      while (!q.empty()) {
                                                                                        int \dot{\mathbf{u}} = \mathbf{q}. \, \text{front}();
     return make_pair(tot_weight, n_matches);
                                                                                         q.pop();
                                                                                         for (edge &e : G[u])
   void add_edge(int u, int v, int w)
   \{g[u][v].w = g[v][u].w = w; \}
                                                                                            if (!~dis[e.to] & e.flow != e.cap)
                                                                                              q.push(e.to), dis[e.to] = dis[u] + 1;
        SW-mincut [6621f5]
4.7
                                                                                      return dis[t] != -1;
int maxflow(int _s, int _t) {
                                                                                      s = \underline{s}, t = \underline{t};

int flow = 0, df;
   \label{eq:man_section}  \begin{array}{ll} \text{int} & \text{vst} \left[ \text{MXN} \right] \,, & \text{edge} \left[ \text{MXN} \right] \left[ \text{MXN} \right] \,, & \text{wei} \left[ \text{MXN} \right] ; \end{array}
                                                                                      while (bfs()) {
   void init(int n)
                                                                                         fill_n(cur, n + 3, 0);
     REP fill_n (edge[i], n, 0);
                                                                                         while ((df = dfs(s, INF))) flow += df;
  void addEdge(int u, int v, int w){
  edge[u][v] += w; edge[v][u] += w;
                                                                                      return flow;
                                                                                   bool solve() {
   int search(int &s, int &t, int n){
                                                                                      int sum = 0;
     fill_n(vst, n, 0), fill_n(wei, n, 0);
                                                                                      for (int i = 0; i < n; ++i)
     s = t = -1;
                                                                                         if (cnt[i] > 0)
     int mx, cur;
                                                                                         \begin{array}{lll} & \text{add\_edge}(n+1,\ i\ ,\ cnt[i])\ ,\ sum\ +=\ cnt[i];\\ & \text{else} & \text{if}\ (cnt[i]<0)\ add\_edge(i\ ,\ n+2,\ -cnt[i])\ ; \end{array}
     for (int j = 0; j < n; +++j) {
        mx = -1, cur = 0;
                                                                                      if (\text{sum }!=\text{maxflow}(n+1, n+2)) \text{ sum }=-1;
        REP if (wei[i] > mx) cur = i, mx = wei[i];
                                                                                      for (int i = 0; i < n; ++i)
        vst[cur] = 1, wei[cur] = -1;
                                                                                         if (cnt[i] > 0)
        s = t; t = cur
                                                                                           G[n + 1].pop_back(), G[i].pop_back();
lse if (cnt[i] < 0)
        REP if (!vst[i]) wei[i] += edge[cur][i];
                                                                                           G[i].pop\_back(), G[n + 2].pop\_back();
     return mx;
                                                                                      return sum != -1;
   int solve(int n) {
                                                                                   int solve(int _s, int
     int res = INF;
                                                                                      add\_edge(\_t, \_s, INF);
     for (int x, y; n > 1; n--){
                                                                                      if (!solve()) return -1; // invalid flow
        res = min(res, search(x, y, n));
REP edge[i][x] = (edge[x][i] += edge[y][i]);
                                                                                      int x = G[\_t].back().flow;
                                                                                      return G[_t].pop_back(), G[_s].pop_back(), x;
        REP {
           edge[y][i] = edge[n - 1][i];
edge[i][y] = edge[i][n - 1];
                                                                                };
        4.9 Gomory Hu tree* [11be99]
                                                                                MaxFlow Dinic;
     return res;
                                                                                int g [MAXN];
} sw;
                                                                                 void GomoryHu(int n) { // 0-base
                                                                                   fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
        BoundedFlow*(Dinic*) [4a793f]
                                                                                      Dinic.reset();
struct BoundedFlow { // 0-base
                                                                                      \begin{array}{lll} add\_edge(i\;,\;g[\,i\,]\;,\;Dinic.maxflow(\,i\;,\;g[\,i\,])\,)\,;\\ for\;(\,int\;\;j=\,i\;+\;1;\;\;j<=\,n;\;+\!\!+\!\!j\,) \end{array}
   struct edge {
     int to, cap, flow, rev;
                                                                                          if (g[j] = g[i] &  \text{ $\sim$ Dinic.dis}[j] ) 
                                                                                           g\,[\,j\,] \;=\; i\;;
  vector<edge> G[N];
int n, s, t, dis[N], cur[N], cnt[N];
```

struct MinCostCirculation { // 0-base
 struct Edge {

4.10 Minimum Cost Circulation* [ba97cf]

```
ll from, to, cap, fcap, flow, cost, rev; } *past[N];
       vector < Edge > G[N];
11 dis [N], inq [N], n;
       void BellmanFord(int s) {
              fill\_n \left( \, dis \, , \, \, n \, , \, \, INF \right) \, , \  \, fill\_n \left( \, inq \, , \, \, n \, , \, \, \, 0 \right) ;
              queue<int> q;
              auto relax = [&](int u, ll d, Edge *e) {
  if (dis[u] > d) {
                            dis[u] = d, past[u] = e;
                            if (!inq[u]) inq[u] = 1, q.push(u);
                    }
              relax(s, 0, 0);
              while (!q.empty()) {
                    int u = q.front();
                     q.pop(), inq[u] = 0;
                     for (auto &e : G[u])
                            if (e.cap > e.flow)
                                  relax(e.to, dis[u] + e.cost, &e);
       }
       void try_edge(Edge &cur) {
              if (cur.cap > cur.flow) return ++cur.cap, void();
              BellmanFord(cur.to);
              if (dis[cur.from] + cur.cost < 0) {
    ++cur.flow, --G[cur.to][cur.rev].flow;</pre>
                     for (int
                          \begin{tabular}{lll} $i=cur.from; &past[i]; &i=past[i]->from) & auto &e=*past[i]; &++e.flow, &--G[e.to][e.rev].flow; & auto &e=*past[i]->from &e=*past[i]-
                    }
             ++cur.cap;
       void solve(int mxlg) {
  for (int b = mxlg; b >= 0; --b) {
                     for (int i = 0; i < n; ++i)
                           for (auto &e : G[i])
                     e cap *= 2, e flow *= 2;
for (int i = 0; i < n; ++i)
for (auto &e : G[i])
                                  if (e.fcap >> b & 1)
                                        try_edge(e);
             }
       void init (int \underline{\phantom{a}}n) { n = \underline{\phantom{a}}n;
              for (int i = 0; i < n; ++i) G[i].clear();
       void add_edge(ll a, ll b, ll cap, ll cost) {
             G[a].pb(Edge
                             \{a, b, 0, cap, 0, cost, SZ(G[b]) + (a = b)\}\);
             G[b].pb(Edge\{b, a, 0, 0, -cost, SZ(G[a]) - 1\});
} mcmf; // O(VE * ElogC)
```

4.11 Flow Models

- $\bullet \quad Maximum/Minimum flow with lower bound/Circulation problem\\$
 - 1. Construct supersource S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite $\operatorname{graph}(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct supersource S and $\sinh T$
 - 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1

- 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) = (0, d(v))
- 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer T
- $2. \ \ {\it Constructa\, max\, flow\, model}, {\it let}\, K\, {\it be\, the \, sum\, of \, all\, weights}$
- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \rightarrow v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - $3. \ \ The {\it mincut} is equivalent to the {\it maximum} profit of a subset of projects.$
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
 - 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \ge 0$$

4.12 matching

- 最大匹配+最小邊覆蓋=V
- 最大獨立集+最小點覆蓋=V
- 最大匹配=最小點覆蓋
- 最小路徑覆蓋數=V-最大匹配數

5 String

5.1 KMP [5a0728]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
  }
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  }
  return ans;
}</pre>
```

5.2 Z-value* [b47c17]

5.3 Manacher* [lad8ef]

```
int z[MAXN]; // 0-base
/* center i: radius z[i * 2 + 1] / 2
    center i, i + 1: radius z[i * 2 + 2] / 2
    both aba, abba have radius 2 */
void Manacher(string tmp) {
    string s = "%";
    int l = 0, r = 0;
    for (char c : tmp) s.pb(c), s.pb("%");
    for (int i = 0; i < SZ(s); ++i) {
        z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
        while (i - z[i] >= 0 && i + z[i] < SZ(s)
        && s[i + z[i]] == s[i - z[i]]) ++z[i];
}</pre>
```

int R = q.front();

q.pop(), ord[t++] = R;

```
\quad \text{for (int } i = 0; \ i < sigma; ++i)
      if (z[i] + i > r) r = z[i] + i, l = i;
                                                                                        if (~nx[R][i])
}
                                                                                           \begin{array}{lll} & \text{int } X = rnx \left[ R \right] \left[ \ i \ \right] = nx \left[ R \right] \left[ \ i \ \right], \ Z = fl \left[ R \right]; \\ & \text{for } (; \ Z \&\& \ ! \sim nx \left[ Z \right] \left[ \ i \ \right]; \ ) \ Z = fl \left[ Z \right]; \\ & fl \left[ X \right] = Z \ ? \ nx \left[ Z \right] \left[ \ i \ \right] \ : \ 1, \ q.push (X); \end{array}
 5.4 SAIS* [6f26bc]
auto sais(const auto &s) {
                                                                                        else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
   const int n = SZ(s), z = ranges :: max(s) + 1;
                                                                                  }
   if (n = 1) return vector{0};
   vector < int > c(z); for (int x : s) + c[x];
                                                                                void solve() {
  for (int i = top - 2; i > 0; --i)
    cnt[fl[ord[i]]] += cnt[ord[i]];
   partial\_sum(ALL(c), begin(c));
   } ac;
      t[i] = (
          s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
                                                                             5.6 Smallest Rotation [4f469f]
   auto is_lms = views::filter([&t](int x) {
      return x && t[x] && !t[x - 1];
                                                                             string mcp(string s) {
                                                                                int n = SZ(s), i = 0, j = 1;
   auto induce = [&] {
                                                                                s += s;
     for (auto x = c; int y : sa)
                                                                                while (i < n \& j < n) {
      if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
for (auto x = c; int y : sa | views::reverse)
if (y--) if (t[y]) sa[--x[s[y]]] = y;
                                                                                   int k = 0;
                                                                                   while (k < n \& s[i + k] = s[j + k]) ++k;
                                                                                   if (s[i + k] \le s[j + k]) j += k + 1;
                                                                                   else i += k + 1;
   vector<int> lms, q(n); lms.reserve(n);
                                                                                   if (i == j) ++j;
   for (auto x = c; int i : I | is_lms)

q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
                                                                                int ans = i < n ? i : j;
   induce(); vector < int > ns(SZ(lms));
                                                                                return s.substr(ans, n);
   for (int j = -1, nz = 0; int i : sa \mid is\_lms) {
     if (j >= 0) {
        5.7 De Bruijn sequence* [a09470]
         ns[q[i]] = nz += lexicographical\_compare(
                                                                             constexpr int MAXC = 10, MAXN = 1e5 + 10;
             \begin{array}{l} \text{begin}(s) + j, & \text{begin}(s) + j + \text{len}, \\ \text{begin}(s) + i, & \text{begin}(s) + i + \text{len}); \end{array}
                                                                             struct DBSeq {
                                                                                int C, N, K, L, buf [MAXC * MAXN]; // K <= \mathbb{C}^N
                                                                                void dfs(int *out, int t, int p, int &ptr) {
      j \ = \ i \ ;
                                                                                   if (ptr >= L) return;
                                                                                   if (t > N) {
if (N \% p) return;
   fill(ALL(sa), 0); auto nsa = sais(ns);
   for (auto x = c; int y : nsa \mid views::reverse)

y = lms[y], sa[--x[s[y]]] = y;
                                                                                      for (int i = 1; i \le p \&\& ptr < L; ++i)
                                                                                        \operatorname{out}[\operatorname{ptr}++] = \operatorname{buf}[i];
   return induce(), sa;
                                                                                   } else
                                                                                     \begin{array}{l} buf[t] = buf[t - p], \ dfs(out, \ t + 1, \ p, \ ptr); \\ for \ (int \ j = buf[t - p] + 1; \ j < C; ++j) \end{array}
 // sa[i]: sa[i]-th suffix
 is the i-th lexicographically smallest suffix. // hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
                                                                                        buf[t] = j, dfs(out, t + 1, t, ptr);
 struct Suffix {
   int n; vector<int> sa, hi, ra;
                                                                                void solve(int _c, int _n, int _k, int *out) {
   Suffix
     \begin{array}{l} (const\ auto\ \&\_s,\ int\ \_n):\ n(\_n)\,,\ hi\,(n)\,,\ ra\,(n)\ \{\\ vector<int>\ s(n+1);\ //\ s\,[n]\ =\ 0;\\ copy\_n(\_s,\ n,\ begin\,(s));\ //\ \_s\ shouldn\ 't\ contain\ 0 \end{array}
                                                                                  int p = 0;
                                                                                  sa = sais(s); sa.erase(sa.begin())
      for (int i = 0; i < n; ++i) ra[sa[i]] = i;
      for (int i = 0, h = 0; i < n; ++i) {
  if (!ra[i]) { h = 0; continue; }
                                                                             } dbs;
                                                                             5.8 Extended SAM* [64c3b7]
         for (int j = sa[ra[i] - 1]; max
                   j) + h < n \& s[i + h] = s[j + h];) + h;
                                                                             hi[ra[i]] = h ? h-- : 0;
   }
};
                                                                                int cnt[N * 2]; // occurence
        Aho-Corasick Automatan* [794a77]
                                                                                int newnode()
                                                                                   fill_n (next[tot], CNUM, 0);
 struct AC_Automatan {
                                                                                   len[tot] = cnt[tot] = link[tot] = 0;
   int nx[len][sigma], fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
                                                                                   return tot++;
                                                                                void init() { tot = 0, newnode(), link[0] = -1; }
   int newnode() {
      fill_n(nx[top], sigma, -1);
                                                                                int insertSAM(int last, int c) {
                                                                                  int cur = next[last][c];
len[cur] = len[last] + 1;
      return top++;
                                                                                   int p = link[last];
   void init() \{ top = 1, newnode(); \}
   int input(string &s) {
                                                                                   while (p != -1 && ! next[p][c])
                                                                                   next[p][c] = cur, p = link[p];
if (p = -1) return link[cur] = 0, cur;
      int X = 1;
      for (char c : s) {    if (!\simnx[X][c - 'A']) nx[X][c - 'A'] = newnode();    X = nx[X][c - 'A'];
                                                                                   int q = next[p][c];
                                                                                   if (len
                                                                                        [p] + 1 = len[q] return link[cur] = q, cur;
      return X; // return the end node of string
                                                                                   int clone = newnode();
                                                                                   for (int i = 0; i < ONUM; ++i)
   void make_fl() {
                                                                                     next[
      queue<int> q;
                                                                                           clone][i] = len[next[q][i]] ? next[q][i] : 0;
      q.push(1), fl[1] = 0;
                                                                                   len[clone] = len[p] + 1;
      for (int t = 0; !q.empty(); ) {
                                                                                   while (p != -1 && next[p][c] == q)
```

next[p][c] = clone, p = link[p];

link[link[cur] = clone] = link[q];

```
link[q] = clone;
                                                                     void main_lorentz(const string &s, int sft = 0) {
     return cur;
                                                                       const int n = s.size();
                                                                       if (n = 1) return;
const int nu = n / 2, nv = n - nu;
   void insert (const string &s) {
     int cur = 0;
                                                                       const string u = s.substr(0, nu), v = s.substr(nu),
     for (auto ch : s) {
                                                                              ru(u.rbegin
       int &nxt = next[cur][int(ch - 'a')];
                                                                                   (), u.rend()), rv(v.rbegin(), v.rend());
                                                                       \begin{array}{ll} main\_lorentz(u, sft), \; main\_lorentz(v, sft + nu); \\ const \; auto \; z1 = Zalgo(ru), \; z2 = Zalgo(v + '\#' + u) \end{array}
        if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
                                                                                   z3 = Zalgo(ru' + '\#' + rv), z4 = Zalgo(v);
     }
                                                                       auto get_z = [](const vector<int>&z, int i) {
   }
   void build() {
                                                                         return
     queue < int > q;
                                                                                (0 \le i \text{ and } i \le (int)z.size()) ? z[i] : 0; };
                                                                       auto add_rep
     q.push(0);
     while (!q.empty()) {
                                                                             = [&](bool left, int c, int l, int k1, int k2) {
       int cur = q.front();
                                                                         const
                                                                               int L = max(1, l - k2), R = min(l - left, k1);
        q.pop();
        for (int i = 0; i < CNUM; ++i)
  if (next[cur][i])</pre>
                                                                          if (L > R) return;
                                                                         if (left)
            q.push(insertSAM(cur, i));
                                                                               rep[l].emplace\_back(sft + c - R, sft + c - L);
                                                                          else rep[1].emplace_back
                                                                              (sft + c - R - l + 1, sft + c - L - l + 1);
     vector < int > lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];
     partial_sum(ALL(lc), lc.begin());
                                                                       for (int cntr = 0; cntr < n; cntr++) {
                                                                         int 1, k1, k2;
     for (int i
          = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
                                                                         if (cntr < nu) {
                                                                            l = nu - cntr;
   void solve() {
                                                                            k1 = get_z(z1, nu - cntr);
     for (int i = tot - 2; i >= 0; --i)
                                                                            k2 = get_z(z2, nv + 1 + cntr);
        cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
                                                                         } else {
                                                                            l = cntr - nu + 1;
 };
                                                                            k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
                                                                            k2 = get_z(z4, (cntr - nu) + 1);
5.9 PalTree* [d7d2cf]
                                                                          if (k1 + k2 >= 1)
struct palindromic_tree {
                                                                            add\_rep\,(\,cntr\,<\,nu\,,\ cntr\,,\ l\,,\ k1\,,\ k2\,)\,;
   struct node {
     int next[26], fail, len;
                                                                    int cnt, num; // cnt: appear times, num: number of
                      // pal. suf.
                                                                     6
                                                                          Math
     node(int \ l = 0)': fail(0), len(1), cnt(0), num(0) {
                                                                     6.1 ax+by=gcd(only exgcd *) [7b833d]
       for (int i = 0; i < 26; ++i) next[i] = 0;
                                                                     pll exgcd(ll a, ll b)
                                                                       if (b = 0) return pll(1, 0);
   vector<node> St;
                                                                       ll p = a / b;
   vector<char> s;
                                                                       pll q = exgcd(b, a \% b);
   int last, n;
                                                                       palindromic\_tree() : St(2), last(1), n(0) 
     St[0].fail = 1, St[1].len = -1, s.pb(-1);
                                                                     /* ax+by=res, let x be minimum non-negative
                                                                    g, p = gcd(a, b), exgcd(a, b) * res / g
   inline void clear() {
                                                                     \begin{array}{l} \text{if } p.X < 0: \ t = (abs(p.X) + b \ / \ g \ - \ 1) \ / \ (b \ / \ g) \\ \text{else: } t = -(p.X \ / \ (b \ / \ g)) \\ p += (b \ / \ g, \ -a \ / \ g) \ * \ t \ */ \\ \end{array} 
     St.clear(), s.clear(), last = 1, n = 0; St.pb(0), St.pb(-1);
     St[0]. fail = 1, s.pb(-1);
                                                                     6.2 Floor and Ceil [692c04]
   int floor (int a, int b)
                                                                    { return \hat{a} / \hat{b} - (a % \hat{b} && (a < 0) \hat{} (b < 0)); } int ceil(int a, int b)
       x = St[x]. fail;
     return x;
                                                                    { return a / b + (a \% b \&\& (a < 0) \cap (b > 0)); }
   inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
  int cur = get_fail(last);
                                                                     6.3 Floor Enumeration [7cbcdf]
                                                                      / enumerating x = floor(n / i), [l,
     if (!St[cur].next[c]) {
  int now = SZ(St);
                                                                    for (int l = 1, r; l <= n; l = r + 1) {
int x = n / l;
        St.pb(St[cur].len + 2);
                                                                       r = n / x;
        St[now].fail =
                                                                    }
          St[get_fail(St[cur].fail)].next[c];
        St[cur].next[c] = now;
                                                                     6.4 Mod Min [9118e1]
        St[now].num = St[St[now].fail].num + 1;
                                                                    // \min\{k \mid 1 \le ((ak) \mod m) \le r\}, no solution -> -1 ll \max_{\min}(ll \ a, \ ll \ m, \ ll \ l, \ ll \ r)  { if (a == 0) return l \ ? \ -1 : \ 0; if (ll \ k = (l + a \ -1) \ / \ a; \ k * a <= r)
     last = St[cur].next[c], ++St[last].cnt;
   inline void count() { // counting cnt
     auto i = St.rbegin();
                                                                         return k;
     for (; i != St.rend(); ++i) {
                                                                       ll\ b = m\ /\ a\,,\ c = m\ \%\ a\,;
       St[i->fail].cnt += i->cnt;
                                                                       return (1 + y * c + a - 1) / a + y * b;
   inline int size() { // The number of diff. pal.
return SZ(St) - 2;
                                                                     6.5 Linear Mod Inverse [5a4cbf]
 };
                                                                    inv[1] = 1:
                                                                     for ( int i = 2; i
 5.10 Main Lorentz [615b8f]
                                                                          \langle = N; ++i \rangle inv[i] = ((mod-mod/i)*inv[mod%i])%mod;
| \text{vector} < \text{pair} < \text{int}, \text{ int} \gg \text{rep}[kN]; // 0-\text{base}[l, r]
```

6.6 Linear Filter Mu [ac2ac3]

6.7 Gaussian integer gcd [763e59]

6.8 GaussElimination [6308be]

```
\begin{tabular}{ll} \begin{tabular}{ll} void $GAS(V<V<double>>&vc) & \{ \end{tabular}
     int len = vc.size();
     for (int i = 0; i < len; ++i)
           int idx = find_if(vc.begin()+i, vc.end(), [&](
               auto&v) {return v[i] != 0;} ) - vc.begin();
           if ( idx = len ) continue
           if(i!=idx) swap(vc[idx], vc[i]);
          double pivot = vc[i][i];
          for_each( vc[i].begin(), vc
                [i].end(), [\&](auto\&a) {a/=pivot;});
           for ( int j = 0; j < len; ++j ) {
    if ( i == j ) continue;
               if( vc[j][i] != 0 ) {
    double mul = vc[j][i]/vc[i][i];
    reconstruction | vec[j][i]/vc[i][i];
                     transform(vc[j].begin(), vc[j].end
                          (), vc[i].begin(), vc[j].begin(),
                               [&](auto &a, auto &b) {
                               return a-b*mul;
                               });
               }
          }
     }
};
```

6.9 floor sum* [49de67]

```
11 floor_sum(ll n, ll m, ll a, ll b) {
    ll ans = 0;
    if (a >= m)
        ans += (n - 1) * n * (a / m) / 2, a %= m;
    if (b >= m)
        ans += n * (b / m), b %= m;
    ll y_max
        = (a * n + b) / m, x_max = (y_max * m - b);
    if (y_max == 0) return ans;
    ans += (n - (x_max + a - 1) / a) * y_max;
    ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
    return ans;
}// sum^{{
    n-1}_0 floor((a * i + b) / m) in log(n + m + a + b)
```

6.10 Miller Rabin* [06308c]

```
while (--t)
  if ((x = mul(x, x, n)) == n - 1) return 1;
return 0;
}
```

6.11 Fraction [4ab37a]

```
struct fraction {
   ll n, d;
   fraction
         (const ll &_n=0, const ll &_d=1): n(_n), d(_d) {
      ll\ t\ =\ \gcd(n,\ d)\,;
      n \neq t, d \neq t;
      if'(d < 0)'n = -n, d = -d;
   fraction operator - () const
   { return fraction(-n, d); }
fraction operator+(const fraction &b) const
{ return fraction(n * b.d + b.n * d, d * b.d); }
   fraction operator-(const fraction &b) const { return fraction(n * b.d - b.n * d, d * b.d); }
   fraction operator*(const fraction &b) const { return fraction(n * b.n, d * b.d); }
   fraction operator/(const fraction &b) const
   { return fraction(n * b.d, d * b.n); }
   void print() {
      cout << n;
       \  \  \, \text{if} \  \  \, (d \ != \ 1) \  \, \text{cout} << \ ^{n}/^{n} << \ ^{d}; \\
};
```

6.12 Simultaneous Equations [a231be]

```
struct matrix { //m variables, n equations
  int n, m;
   fraction M[MAXN][MAXN + 1], sol [MAXN];
   int solve() { //-1: inconsistent, >= 0: rank
     for (int^{'}i = 0; i < n; ++i) {
      int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;
       if (piv == m) continue
       for (int j = 0; j < n; +++j) {
         if (i == j) continue
         fraction \ tmp = -M[\,j\,][\,piv\,] \ / \ M[\,i\,][\,piv\,]\,;
         for (int k = 0; k < 0
              m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
      }
    int rank = 0;
    for (int i = 0; i < n; ++i) {
      int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;
       if (piv = m \&\& M[i][m].n) return -1;
       else if (piv
            < m) ++rank, sol[piv] = M[i][m] / M[i][piv];
     return rank;
  }
};
```

6.13 Pollard Rho* [fdef9b]

6.14 Simplex Algorithm [6b4566]

```
 \begin{array}{l} \textbf{const} \;\; \textbf{int} \;\; \textbf{MAXN} = 11000, \; \textbf{MAXM} = 405; \\ \textbf{const} \;\; \textbf{double} \;\; \textbf{eps} = 1\text{E-}10; \\ \textbf{double} \;\; \textbf{a} \; [\textbf{MAXN}] \; [\textbf{MAXM}] \;, \;\; \textbf{b} \; [\textbf{MAXN}] \;, \;\; \textbf{c} \; [\textbf{MAXM}]; \\ \textbf{double} \;\; \textbf{d} \; [\textbf{MAXN}] \; [\textbf{MAXM}] \;, \;\; \textbf{x} \; [\textbf{MAXM}] \;; \\ \end{array}
```

```
\begin{array}{lll} & \text{int ix} \left[ \text{MAXN} + \text{MAXM} \right]; \ // \ !!! \ \text{array all indexed from 0} \\ // \ \text{max} \{ \text{cx} \} \ \text{subject to} \ \{ \text{Ax}\!\!=\!\!b, x \!\!>\!\!=\!\!0 \} \end{array}
// n: constraints, m: vars !!!
    x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  ++m;
   fill\_n\,(d\,[\,n\,]\;,\;m\,+\,1\,,\;\,0)\,;
   fill_n(d[n+1], m+1, 0);

iota(ix, ix + n + m, 0);
   int r = n, s = m - 1;
   for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];
      d[i][m - 1] = 1;
      d[i][m] = b[i];
      if (d[r][m] > d[i][m]) r = i;
   copy_n(c, m - 1, d[n]);
d[n + 1][m - 1] = -1;
   for (double dd;; ) {
      if (r < n)
         swap(ix[s], ix[r + m]);
d[r][s] = 1.0 / d[r][s];
for (int j = 0; j <= m; ++j)
    if (j != s) d[r][j] *= -d[r][s];</pre>
          for (int i = 0; i <= n + 1; ++i) if (i != r) {
  for (int j = 0; j <= m; ++j) if (j != s)
  d[i][j] += d[r][<math>j] * d[i][s];
  d[i][s] *= d[r][s];
          }
      }
      r = s = -1;
      for (int j = 0; j < m; ++j)

if (s < 0 || ix[s] > ix[j]) {

if (d[n + 1][j] > eps ||

(d[n + 1][j] > -eps && d[n][j] > eps))
       if (s < 0) break;
       for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
          if (r < 0 | |
                 (dd = d[r][m]
                          / d[r][s] - d[i][m] / d[i][s]) < -eps ||
                 (dd < eps & ix[r + m] > ix[i + m]))
             r = i;
      if (r < 0) return -1; // not bounded
   if (d[n + 1][m] < -eps) return -1; // not executable
   double ans = 0;
   fill_n(x, m, 0);
   for (int i = m; i <
      \begin{array}{ll} (n \cdot i - i \cdot i, i - i \cdot i) & \text{if } (i \cdot i) = 0 \\ \text{if } (i \cdot i) & \text{if } (i \cdot i) = 0 \\ \text{ans } += d[i - m][m] & \text{c}[i \cdot i]; \\ \end{array}
          x[ix[i]] = d[i-m][m];
   return ans;
```

6.14.1 Construction

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c$

 $\overline{\mathbf{x}}$ and $\overline{\mathbf{y}}$ are optimalified only if for all $i \in [1, n]$, either $\overline{x}_i = 0$ or $\sum_{j=1}^m A_{ji} \overline{y}_j = c_i$ holds and for all $i \in [1,m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij} \bar{x}_j = b_j$ holds.

1. In case of minimization, let $c'_i = -c_i$

- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3. $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
 - $\bullet \quad \sum_{1 \le i \le n}^{-} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.15 chineseRemainder [a53b6d]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
  \begin{array}{l} ll \;\; g = \gcd{(m1,\; m2)}\,; \\ if \;\; ((x2 \; - \; x1) \;\% \; g) \;\; return \;\; \text{-1}; \; \text{// no sol} \end{array}
  m1 /= g; m2 /= g;
   pll p = exgcd(m1, m2);
  ll lcm = m1 * m2 * g;

ll res = p.first * (x^2 - x^1) * m1 + x1;
   // be careful with overflow
```

```
15
   return (res % lcm + lcm) % lcm;
6.16 Factorial without prime factor* [c324f3]
  / O(p^k + \log^2 n), pk = p^k
 11 prod [MAXP];
 ll fac_no_p(ll n, ll p, ll pk) {
   \operatorname{prod}[0] = 1;
   for (int i = 1; i <= pk; ++i)
if (i % p) prod[i] = prod[i - 1] * i % pk;
      else prod[i] = prod[i - 1];
   11 \text{ rt} = 1;
   for (; n; n /= p) {
     rt = rt * mpow(prod[pk], n / pk, pk) % pk;
     rt = rt * prod[n \% pk] \% pk;
   return rt;
} // (n! without factor p) % p^k
6.17 PiCount* [cad6d4]
ll PrimeCount(ll n) { // n ~ 10^13 \Rightarrow < 2s
   if (n \le 1) return 0;
   int v = sqrt(n), s = (v + 1) / 2, pc = 0;
   vector < int > smalls(v + 1), skip(v + 1), roughs(s);
   vector<ll> larges(s);
   for (int i = 2; i \le v; ++i) smalls[i] = (i + 1) / 2;
   for (int i = 2, i < v, i + 1) smalls[i] = for (int i = 0; i < s; ++i) { roughs[i] = 2 * i + 1; larges[i] = (n / (2 * i + 1) + 1) / 2;
   for (int p = 3; p <= v; ++p) {
  if (smalls[p] > smalls[p - 1]) {
   int q = p * p;
        ++pc;
        if (1LL * q * q > n) break;
        skip[p] = 1;
        for (int i = q; i \le v; i += 2 * p) skip[i] = 1;
        int ns = 0;
        for (int k = 0; k < s; ++k) {
          int i = roughs[k];
          if (skip[i]) continue;

ll d = lLL * i * p;

larges[ns] = larges[k] - (d <= v ? larges
               [smalls[d] - pc] : smalls[n / d]) + pc;
          roughs [ns++] = i;
        }
        s = ns;
        for (int j = v / p; j >= p; --j) {
          int c =
          }
     }
   for (int k = 1; k < s; ++k) {
     const ll m = n / roughs[k];
ll t = larges[k] - (pc + k - 1);
      for (int l = 1; l < k; ++l) {
       int p = roughs[1];
        if (1LL * p * p > m) break;
t -= smalls[m / p] - (pc + l - 1);
     larges[0] = t;
   }
   return larges [0];
6.18 Discrete Log* [da27bf]
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {
```

```
\begin{array}{l} p\,[\,y\,] \;=\; i\;;\\ y \;=\; 1LL\; *\; y\; *\; x\; \%\; m;\\ b \;=\; 1LL\; *\; b\; *\; x\; \%\; m; \end{array}
    for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;
        if (p.find(s) != p.end()) return i + kStep - p[s];
    return -1;
int DiscreteLog(int x, int y, int m) {
```

```
if (m = 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
   if (s == y) return i;
s = 1LL * s * x % m;
  if (s = y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
6.19 Berlekamp Massey [3eb6fa]
template <typename T>
```

```
vector<T> BerlekampMassey(const vector<T> &output) {
     vector \langle I \rangle degree amplifies by (const. vector \langle I \rangle degree vector \langle I \rangle d(SZ(output) + 1), me, he; for (int f = 0, i = 1; i <= SZ(output); ++i) { for (int j = 0; j < SZ(me); ++j) d[i] += output[i - j - 2] * me[j]; if ((d[i] -= output[i - 1]) == 0) continue;
            if (me.empty()) {
                  me. resize (f = i);
                  continue;
          f
vector<T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
for (T x : he) o.pb(x * k);
o.resize(max(SZ(o), SZ(me)));
for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
           me = o;
```

6.20 Primes

return me;

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 \ 98789101 \ 987777733 \ 999991921 \ 1010101333
     1010102101 \ 1000000000039 \ 100000000000037
     2305843009213693951 \ \ 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

Theorem

```
Lucas's Theorem
```

For non-negative integer n, m and prime P, $C(m,n) \operatorname{mod} P = C(m/M,n/M) * C(m\%M,n\%M) \operatorname{mod} P$ = mult_i(C(m_i,n_i)) $where \verb|m_i| is the i-th digit of \verb|m| in base P.$

Kirchhoff's theorem

 $A_{\{ii\}} = deg(i), A_{\{ij\}} = (i,j) \in ?-1:0$ Deleting any one row, one column, and calthedet (A)

$Nth\,Catalan\,recursive\,function:$

```
C_0=1, C_{n+1}=C_n*2(2n+1)/(n+2)
```

```
Mobius Formula
u(n)=1, if n=1
 (-1)<sup>m</sup>,若n無平方數因數,且n=p1*p2*p3*...*pk
0 ,若n有大於1的平方數因數
-Property
1.(積性函數)u(a)u(b)=u(ab)
```

2. $_{\{d|n\}} u(d) = [n == 1]$

```
Mobius Inversion Formula
\begin{array}{l} \mathrm{if} \ f(n) = \ \_\{d|n\}\,g(d) \\ \mathrm{then} \ g(n) = \ \_\{d|n\}\,u(n/d)f(d) \\ = \ \_\{d|n\}\,u(d)f(n/d) \end{array}
```

-Application

the number/power of gcd(i, j) = k-Trick

分塊,O(sqrt(n))

Chinese Remainder Theorem (m_i 兩兩互質)

```
x=a_1 \pmod{m_1}
x=a_2 \pmod{m_2}
x=a_i \pmod{m_i}
construct a solution:
Let M=m 1*m 2*m 3*...*m n
\mathrm{Let}\, \mathrm{M}\_\mathrm{i}\!=\!\mathrm{M}/\mathrm{m}\_\mathrm{i}
```

```
t_i = 1/M_i
  t_i^*M_i=1 \pmod{m_i}
 solutionx = a_1*t_1*M_1 + a_2*t_2*M_2 + ... + a_n*t_n*M_n + k*M_n + 
  =k*M+ a_i*t_i*M_i, kispositive integer.
 under mod M, there is one solution x = a\_i*t\_i*M\_i
Burnside's lemma
|G|*|X/G| = sum(|X^g|) where gin G
 總方法數:每一種旋轉下不動點的個數總和除以旋轉的方法數
```

6.22 Estimation

 $n \hspace{.2cm} | 2345 \hspace{.05cm} 6 \hspace{.05cm} 7 \hspace{.05cm} 8 \hspace{.05cm} 9 \hspace{.05cm} 20 \hspace{.05cm} 30 \hspace{.05cm} 40 \hspace{.05cm} 50 \hspace{.05cm} 100$ p(n) | 23571115223062756044e42e52e8n |1001e31e6 1e9 1e12 1e15 1e18 $\overline{d(i)}$ 12 32 240 1344 6720 26880 10 3680 $n \mid 23 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \ 13$ B_n 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7

6.23 General Purpose Numbers

• Bernoulli numbers
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$
Strike the first first of the limit Positive of the limit Po

• Stirling numbers of the second kind Partitions of n distinct elements into

exactly kilothes.
$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$x^{k} = \sum_{i=0}^{k} S(n,i)(x)_{i}$$
• Pentagonal number theorem
$$\prod_{n=1}^{\infty} (1-x^{n}) = 1 + \sum_{k=1}^{\infty} (-1)^{k} \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
• Catalan numbers
$$C_{n}^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^{k}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. kj:ss.t. $\pi(j) > \pi(j+1), k+1j$:ss.t. $\pi(j) \ge j$, kj:ss.t. $\pi(j) > j$.

E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

6.24 Tips for Generating Functions

• Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$

```
-A(rx) \Rightarrow r^n a_n
-A(x)+B(x) \Rightarrow a_n+b_n
-A(x)B(x)\Rightarrow\Sigma_{i=0}^n a_i b_{n-i}
- A(x)^{k} \Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n} a_{i_{1}} a_{i_{2}} \dots a_{i_{k}}
-xA(x)' \Rightarrow na_n
-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i
```

• Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$

```
-A(x)+B(x) \Rightarrow a_n+b_n
-A^{(k)}(x) \Rightarrow a_{n+k}
-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}
-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}
```

• Special Generating Function $- (1+x)^n = \sum_{i \ge 0} {n \choose i} x^i$

 $-xA(x) \Rightarrow na_n$

$- \frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{i}{n-1} x^i$ Polynomial

7.1 Fast Fourier Transform [5e2ea2]

```
const int maxn = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = a\cos(-1);
cplx omega[maxn + 1];
\begin{array}{c} \textbf{void} \ \ \textbf{prefft} \ () \ \ \{ \\ \ \ \ \textbf{for} \ \ (\textbf{int} \ \ \textbf{i} \ = \ 0; \ \ \textbf{i} <= \ \text{maxn} \end{array}
              ; ++i) omega[i] = exp(i * 2 * pi / maxn * I);
}
```

```
void bin(vector<cplx> &a, int n) {
        int lg;
                                                                                                                   }
        for (lg = 0; (1 \ll lg) < n; ++lg); --lg;
        vector < cplx > tmp(n);
                                                                                                                void InverseTransform(vector<int> &v, int n) {
                                                                                                                  Transform(v, n);
for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
const int kInv = fpow(n, kMod - 2);
        for (int i = 0; i < n; ++i) {
                int to = 0;
               for (int j = 0; (1 << j) <
                       n; ++j) to |= (((i >> j) \& 1) << (lg - j));
                                                                                                                   for (int i
               tmp[to] = a[i];
                                                                                                                             = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;
       for (int i = 0; i < n; ++i) a[i] = tmp[i];
                                                                                                               7.3 Fast Walsh Transform* [c9cdb6]
                                                                                                               /* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
 \begin{tabular}{ll} \be
       bin(a, n);
                                                                                                               xor: (x, y = (x + y) * op, (x - y) * invop: or, and, xor = -1, -1, 1/2 */
        for (int step = 2; step \leq n; step \leq 1) {
                int to = step \gg 1;
                                                                                                               void fwt(int *a, int n, int op) { //or
                for (int i = 0; i < n; i += step) {
                                                                                                                   for (int L = 2; L \le n; L \le 1)
                       for (int k = 0; k < to; ++k) {
                                                                                                                       for (int i = 0; i < n; i += L)
                              \hat{c}plx x = a[i]
                                                                                                                          for (int j = i; j < i + (L >> 1); ++j)

a[j + (L >> 1)] += a[j] * op;
                              + \text{ to } + \text{ k}] * omega[maxn / step * k];
a[i + to + k] = a[i + k] - x;
                              a[i + k] += x;
                                                                                                               const int N = 21;
                                                                                                               int f
               }
                                                                                                                      N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
       }
                                                                                                                void
}
                                                                                                                        subset_convolution(int *a, int *b, int *c, int L) {
                                                                                                                    // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
void ifft (vector < cplx > &a, int n) {
                                                                                                                   int n = 1 \ll L;
        fft(a, n);
                                                                                                                   for (int i = 1; i < n; ++i)
        reverse(a.begin() + 1, a.end())
                                                                                                                    ct[i] = ct[i & (i - 1)] + 1;
for (int i = 0; i < n; ++i)
        for (int i = 0; i < n; ++i) a[i] /= n;
                                                                                                                      \begin{array}{l} f\left[ct\left[i\right]\right]\left[i\right] = a\left[i\right], \ g\left[ct\left[i\right]\right]\left[i\right] = b\left[i\right]; \\ cr\left(int\ i = 0;\ i <= L;\ +\!\!+\!\!i) \\ fwt(f\left[i\right],\ n,\ 1), \ fwt(g\left[i\right],\ n,\ 1); \end{array}
                                                                                                                    for (int
vector<int> multiply(const vector<
        int > \&a, const vector < int > \&b, bool trim = false) {
                                                                                                                    for (int i = 0; i \leftarrow L; ++i)
        int d = 1;
                                                                                                                       for (int j = 0; j \le i; ++j)
        while
                                                                                                                   for (int x = 0; x < n; ++x)

h[i][x] += f[j][x] * g[i - j][x];

for (int i = 0; i <= L; ++i)
                (d < max(a.size(), b.size())) d <<= 1; d <<= 1;
        vector < cplx > pa(d), pb(d);
        for (int i
                                                                                                                   fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
                 = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
                                                                                                                       c[i] = h[ct[i]][i];
                 = 0; i < b.size(); ++i) pb[i] = cplx(b[i], 0);
        fft(pa, d); fft(pb, d);
        for (int i = 0; i < d; ++i) pa[i] *= pb[i];
                                                                                                               7.4 Polynomial Operation [869cb1]
        ifft (pa, d);
        vector < int > r(d);
        for (int
                                                                                                               \begin{array}{c} \text{fi}\,(s,\,n) \;\; \text{for} \;\; (\text{int}\;i=(\text{int})(s);\;i<(\text{int})(n);\; +\!\!+\!\!i)\\ \text{template}<\text{int}\;\; \text{MAXN},\;ll\;\;P,\;\;ll\;\; RT>\;//\;\; \text{MAXN}=2^k\\ \text{struct}\;\; \text{Poly}\;\;:\;\; \text{vector}<ll>\;\{\;\;//\;\; \text{coefficients}\;\; \text{in}\;\;[0\;,\;P) \end{array}
                 i = 0; i < d; ++i) r[i] = round(pa[i].real());
        if (trim)
                while (r.size() \& r.back() = 0) r.pop\_back();
                                                                                                                   using vector<ll>>::vector;
       return r;
                                                                                                                    static NTT<MAXN, P, RT> ntt;
}
                                                                                                                   int n() const { return (int) size(); } // n() >= 1
Poly(const Poly &p, int m) : vector<ll>(m) {
   copy_n(p.data(), min(p.n(), m), data());
  Prime
                      Root
                                  Prime
                                                        Root
  7681
                      17
                                   167772161
  12289
                      11
                                   104857601
                                                        3
                                                                                                                   Poly& irev()
  40961
                                   985661441
                      3
                                                        3
                                                                                                                            \{ return reverse(data(), data() + n()), *this; \}
  65537
                      3
                                   998244353
                                                                                                                   Poly& isz(int m) { return resize(m), *this; }
Poly& iadd(const Poly &rhs) { // n() = rhs.n()
                      10
                                   1107296257
  786433
                                                        10
  5767169
                      3
                                   2013265921
                                                        31
                                                                                                                       fi(0, n()) if
  (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
  7340033
                      3
                                  2810183681
                                                        11
  23068673
                                   2885681153
                      3
                                                        3
  469762049
                                   605028353
7.2 Number Theory Transform* [7d51db]
                                                                                                                   Poly& imul(ll k) {
                                                                                                                       fi\,(\,0\,,\,\,n\,(\,)\,)\ (\,{}^*t\,h\,i\,s\,)\,[\,i\,]\,\,=\,(\,{}^*t\,h\,i\,s\,)\,[\,i\,]\,\,\,{}^*k\,\,\%\,\,P\,;
vector<int> omega;
void Init() {
                                                                                                                       return *this;
   omega.resize(kN + 1);
   long long x = \text{fpow}(k\text{Root}, (\text{Mod} - 1) / k\text{N});
                                                                                                                   Poly Mul(const Poly &rhs) const {
   omega[0] = 1;
                                                                                                                       int m = 1:
    for (int i = 1; i \le kN; ++i)
                                                                                                                       while (m < n() + rhs.n() - 1) m <<= 1;
       omega[i] = 1LL * omega[i - 1] * x % kMod;
                                                                                                                       Poly X(*this, m), Y(rhs, m);
                                                                                                                       ntt(X.data(), m), ntt(Y.data(), m);
                                                                                                                       fi(0, m) X[i] = X[i] * Y[i] % P;

ntt(X.data(), m, true);
void Transform(vector<int> &v, int n) {
   BitReverse(v, n);
for (int s = 2; s \le n; s \le 1) {
                                                                                                                       return X.isz(n() + rhs.n() - 1);
       int z = s \gg 1;
                                                                                                                   Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
       for (int i = 0; i < n; i += s) {
                                                                                                                       if (n() = 1) return \{ntt.minv((*this)[0])\};
           for (int k = 0; k < z; ++k) {
                                                                                                                       int m = 1;
               int x = 1LL
                                                                                                                       while (m < n() * 2) m <<= 1;
                                                                                                                       Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
                         * v[i + k + z] * omega[kN / s * k] % kMod;
                                                                                                                       Poly Y(*this, m);
                v[i + k + z] = (v[i + k] + kMod - x) \% kMod;
               (v[i + k] + x) \approx kMod;
                                                                                                                       ntt(Xi.data(), m), ntt(Y.data(), m);
                                                                                                                       fi(0, m) {
```

```
Xi[i] *= (2 - Xi[i] * Y[i]) % P;
if ((Xi[i] \%= P) < 0) Xi[i] += P;
   ntt(Xi.data(), m, true);
  return Xi.isz(n());
Poly Sqrt()
       const \{ // Jacobi((*this)[0], P) = 1, 1e5/235ms \}
   if (n()
        = 1) return {QuadraticResidue((*this)[0], P)};
  Poly
        X = Poly(*this \,, \ (n() \,+\, 1) \,\,/\,\, 2).\, Sqrt()\,.\, isz\,(n())\,;
         X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
pair<Poly, Poly> DivMod
(const Poly &rhs) const { // (rhs.)back() != 0
   if (n() < rhs.n()) return \{(0), *this\}; const int m = n() - rhs.n() + 1;
   Poly X(rhs); X. irev(). isz(m);
  Poly Y(*this); Y. irev(). isz(m);
  \begin{array}{l} \mbox{Poly } Q = \mbox{Y.Mul}(\mbox{X.Inv}(\mbox{)}).\mbox{isz}(\mbox{m}).\mbox{irev}(\mbox{)}; \\ \mbox{X = rhs.Mul}(\mbox{Q}), \mbox{Y = *this}; \\ \mbox{fi}(\mbox{0}, \mbox{n}(\mbox{)}) \mbox{ if } ((\mbox{Y}[\mbox{i}] - \mbox{X}[\mbox{i}]) < 0) \mbox{Y}[\mbox{i}] += P; \end{array}
  return {Q, Y. isz (max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
   fi(0,
       ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
   return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
   fi(0, n())
         ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
   return ret;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
   \begin{array}{lll} \textbf{return} & Poly(Y.\,data() \,\, \dot{+} \,\, \dot{n}() \,\, & -1,\,\, Y.\,data() \,\, +Y.\,n())\,; \end{array} 
vector<ll> eval(const
       vector<ll> &x, const vector<Poly> &up) const {
   const int m = (int)x.size();
  if (!m) return { };
   vector<Poly> down(m * 2)
   // \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};
  // fi(2, m *
2) down[i] = down[i / 2].DivMod(up[i]).second;
  down[1] = Poly(up[1])
  vector<ll> y(m);
   fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _treel(const vector<ll> &x) {
  const int m = (int)x.size();
   {\tt vector}{<}{\tt Poly}{\tt > up(m * 2);}
  return up;
     <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate (const vector
     \langle ll \rangle \&x, const vector\langle ll \rangle \&y \rangle  { // 1e5, 1.4s
   const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
fi(0, m) down[m + i] = {z[i]};
   for (int i = m -
        1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
  return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const \{ // (*this)[0] = 0, 1e5/360ms \}
```

```
fi(0, n())
          if'((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;
    return X.Mul(Y).isz(n());
   V/M := P(P - 1). If k >= M, k := k \% M + M.
  Poly Pow(ll k) const {
    int nz = 0;
    while (nz < n() \&\& !(*this)[nz]) ++nz;
if (nz * min(k, (ll)n()) >= n()) return Poly(n());
if (!k) return Poly(Poly \{1\}, n());
    Poly \ X(\ data(\ ) \ + \ nz \ , \ \ data(\ ) \ + \ nz \ + \ n(\ ) \ \ - \ nz \ * \ k) \ ;
     const ll c = ntt.mpow(X[0], k \% (P - 1));
    return X.Ln().imul
         (k \ \% \ P) \ . Exp() \ . imul(c) \ . irev() \ . isz(n()) \ . irev();
  static 11
       LinearRecursion(const vector<11> &a, const vector
       <11> &coef , ll n) { // a_n = \sum c_j a_(n-j)
     const int k = (int)a.size();
     assert((int)coef.size() = k + 1);
    Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\}; fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
    C[k] = 1;
     while (n) {
       if (n \% 2) W = W. Mul(M). DivMod(C). second;
       n \neq 2, M = M.Mul(M).DivMod(C).second;
     11 \text{ ret} = 0;
    fi(0, k) ret = (ret + W[i] * a[i]) \% P;
    return ret;
  }
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template  decltype(Poly_t::ntt) Poly_t::ntt = {};
       Value Polynomial [96cde9]
```

```
struct Poly {
  mint base; // f(x) = poly[x - base]
   vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) \{\}
  mint get_val(const mint &x) (
     if (x \ge base & x < base + SZ(poly))
        return poly[x - base];
     mint rt = 0:
     vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
     for (int i = 1; i < SZ(poly); +i)

lmul[i] = lmul[i - 1] * (x - (base + i - 1));

for (int i = SZ(poly) - 2; i >= 0; -i)

rmul[i] = rmul[i + 1] * (x - (base + i + 1));
     for (int i = 0; i < SZ(poly); ++i)
rt += poly[i] * ifac[i] * inegfac
[SZ(poly) - 1 - i] * lmul[i] * rmul[i];
     return rt:
  void raise() \{ // g(x) = sigma\{base:x\} f(x) \}
     if (SZ(poly) = 1 \& poly[0] = 0)
     mint nw = get\_val(base + SZ(poly));
     poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)
        poly[i] += poly[i - 1];
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

 $\text{for }\beta \text{ being some constant. Polynomial }P \text{ such that }F(P) = 0 \text{ can be found it-}$ eratively. Denoteby Q_k the polynomial such that $F(Q_k) = 0 \pmod{x^{2^k}}$, then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

Geometry

Basic [e38806] 8.1

```
bool same
    (double a, double b) { return abs(a - b) < eps; }
```

yr[m] = max(yr[m], yr[rc[m]]);

```
return m:
struct P {
  double x,
                                                                                bool bound(const point &q, int o, long long d) {
  P() : x(0), y(0) \{ \}
                                                                                   double ds = sqrt(d + 1.0);
 P(): x(0), y(0) {}
P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }
double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
double abe() { return hypot(x y): }
                                                                                   if (q.x < xl[o] - ds | | q.x > xr[o] + ds | |
                                                                                        q\,.\,y\,<\,
                                                                                              yl[o] - ds \mid \mid q.y > yr[o] + ds) return false;
                                                                                   return true;
                                                                                double abs() { return hypot(x, y); }
P unit() { return *this / abs(); }
  P rot(double o) {
                                                                                void dfs (
     double c = cos(o), s = sin(o);
                                                                                      const point &q, long long &d, int o, int dep = 0) {
     return P(c * x - s * y, s * x + c * y);
                                                                                   if (!bound(q, o, d)) return;
                                                                                   long long cd = dist(p[o], q);
  double angle() { return atan2(y, x); }
                                                                                   if (cd != 0) d = min(d, cd);
                                                                                   if ((dep & 1)
                                                                                          && q.x < p[o].x \mid \mid !(dep \& 1) && q.y < p[o].y) {
                                                                                      if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
struct L {
  // ax + by + c = 0
  double a, b, c, o;
                                                                                   } else {
 if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                                      if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                                   }
 P project(P p) { return pa + (pb - pa).unit
    () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs() * (pb - pa).abs()); }
                                                                                void init(const vector<point> &v) {
                                                                                   for (int i = 0; i < v. size(); ++i) p[i] = v[i];
                                                                                   root = build(0, v.size());
                                                                                long long nearest (const point &q) {
  bool inside (
                                                                                   long long res = 1e18;
        P p) \{return min(pa.x, pb.x) <= p.x & p.x <= max \}
                                                                                   dfs(q, res, root);
        return res;
};
                                                                                8.3 Sector Area [c41fb7]
bool SegmentIntersect (P p1, P p2, P p3, P p4) {
  if \ (\max(p1.x\,,\ p2.x)\,<\,\min(p3.x\,,\ p4.x)\ |\ |
                                                                                   calc area of sector which include a, b
                                                                                double SectorArea(Pa, Pb, double r) {
         \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
                                                                                  double o = atan2(a.y, a.x) - atan2(b.y, b.x);

while (o <= 0) o += 2 * pi;

while (o >= 2 * pi) o -= 2 * pi;

o = min(o, 2 * pi - o);

return r * r * o / 2;
  if (\max(p1.y, p2.y) < \min(p3.y, p4.y) \mid |
         \max(p3.y, p4.y) < \min(p1.y, p2.y)) return false;
  return sign((p3 - p1) ^{\circ} (p4 - p1)) * sign((p3 - p2) ^{\circ} (p4 - p2)) <= 0 && sign((p1 - p3) ^{\circ}
            (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
                                                                                8.4 Half Plane Intersection [f7274e]
bool parallel
                                                                                bool jizz (L l1, L l2, L l3) {
     (L x, L y) \{ return same(x.a * y.b, x.b * y.a); \}
                                                                                  P = Intersect(12, 13);
                                                                                   return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
     (L x, L y) { return P(-x.b * y.c + x.c * y.b, x .a * y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
                                                                                bool cmp(const L &a, const L &b){
8.2 KD Tree [375ca2]
                                                                                   return same(
                                                                                        a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
namespace kdt {
int root, lc [maxn],
       \texttt{rc}\left[\texttt{maxn}\right], \ \ \texttt{xl}\left[\texttt{maxn}\right], \ \ \texttt{xr}\left[\texttt{maxn}\right], \ \ \texttt{yl}\left[\texttt{maxn}\right], \ \ \texttt{yr}\left[\texttt{maxn}\right];
                                                                                // availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
point p[maxn];
                                                                                vector<P> HPI(vector<L> &ls){
int build(int l, int r, int dep = 0) {
                                                                                   sort(ls.begin(),ls.end(),cmp);
  if (l = r) return -1;
                                                                                   vector < L > pls(1, ls[0]);
  function < bool (const point &, const point
                                                                                   for (int i=0; i < (int) ls. size ();++i) if (!
     &>> f = [dep](const point &a, const point &b) { if (dep & 1) return a.x < b.x;
                                                                                        same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
                                                                                   deque < int > dq; dq.push_back(0); dq.push_back(1);
     else return a.y < b.y;
                                                                                #define meow(a,b,c
                                                                                        while (dq. size ()>1u & jizz (pls [a], pls [b], pls [c]))
  int m = (1 + r) >> 1;
                                                                                   for (int i=2; i<(int) pls. size ();++i) {
  \begin{array}{l} meow(i\,,dq\,.back\,()\,,dq\,[dq\,.\,size\,()\,-2])dq\,.pop\_back\,()\,;\\ meow(i\,,dq\,[\,0\,]\,,dq\,[\,1\,])dq\,.pop\_front\,()\,; \end{array}
  yl[m] = yr[m] = p[m].y;
                                                                                     dq.push_back(i);
  lc[m] = build(1, m, dep + 1);
  if (~lc[m]) {
      \begin{array}{l} xl\left[m\right] = \min\left(xl\left[m\right], \ xl\left[lc\left[m\right]\right]\right); \\ xr\left[m\right] = \max\left(xr\left[m\right], \ xr\left[lc\left[m\right]\right]\right); \end{array} 
                                                                                  meow (dq
                                                                                         . front(), dq.back(), dq[dq.size()-2])dq.pop\_back();
                                                                                   meow(dq.back(),dq[0],dq[1])dq.pop\_front();
     yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
                                                                                   if (dq.size()<3u) return vector
                                                                                         <P>(); // no solution or solution is not a convex
                                                                                   vector<P> rt;
  rc[m] = build(m + 1, r, dep + 1);
                                                                                   for (int i=0; i<(int)dq. size();++i)rt.push_back
  if (~rc[m]) {
                                                                                         (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
     xl[m] = min(xl[m], xl[rc[m]]);
     xr[m] = max(xr[m], xr[rc[m]]);
     yl[m] = min(yl[m], yl[rc[m]]);
```

8.5 Rotating Sweep Line [0411f0]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
  int n=int(ps.size())
  vector < int > id(n), pos(n);
  vector<pair<int, int>>> line(n*(n-1)/2);
  int m=-1;
  for(int i=0;i< n;++i)for
       (int j=i+1;j< n;++j)line[++m]=make\_pair(i,j); ++m;
  sort(line.begin(),line.end(),[&](const
        pair<int, int> &a, const pair<int, int> &b)->bool{
       if (ps
           [a.first].first==ps[a.second].first)return 0;
       if (ps
           [b.first].first=ps[b.second].first)return 1;
       return (double
           )(ps[a.first].second-ps[a.second].second)/(ps
           [a.first].first-ps[a.second].first) < (double
           ) (ps[b.first].second-ps[b.second].second
           )/(ps[b.first].first-ps[b.second].first);
       });
  for (int i=0; i< n; ++i) id [i]=i;
  sort(id.begin(),id.end(),[&](const
    int &a,const int &b){ return ps[a]<ps[b]; });</pre>
  for (int i=0; i< n; ++i) pos [id[i]]=i;
  for (int i=0; i < m++i)
    auto l=line[i];
    // meow
    tie (pos[l.first],pos[l.second],
         id [pos[l.first]], id [pos[l.second]])=make_tuple
         (pos[l.second],pos[l.first],l.second,l.first);
  }
}
```

Triangle Center [4e8ee9] 8.6

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
 Point res;
 double a1 = atan2(b.y - a.y, b.x - a.x) + pi /
 double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
 double ax = (a.x + b.x) /
 double ay = (a.y + b.y)
 double bx = (c.x + b.x) / 2;
 return Point(ax + r1^* cos(a1), ay + r1^* sin(a1));
Point TriangleMassCenter(Point a, Point b, Point c) {
 return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
 return TriangleMassCenter(a, b
     , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
 Point res;
 double la = len(b - c);
 double lb = len(a - c);
double lc = len(a - b);
 res.x = (
     la *`a.x + lb * b.x + lc * c.x) / (la + lb + lc);
 res.y = (
     la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
 return res;
```

Polygon Center [ee6ff0]

```
Point \ BaryCenter(vector {<\!Point\!> \&\!p\,,\ int\ n})\ \{
 Point res(0, 0);
 res.x += (p[0].x + p[i].x + p[i + 1].x)
   res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
 res.x /= (3 * s);
res.y /= (3 * s);
 return res;
```

Maximum Triangle [3a6d38] 8.8

```
20
double ConvexHullMaxTriangleArea
     (Point p[], int res[], int chnum) {
  double area = 0, tmp;
  res[chnum] = res[0];
  for (int i = 0, j = 1, k = 2; i < \text{chnum}; i++) {
     while (fabs(Cross(p[
          \begin{array}{c} res[j]] - p[res[i]], \ p[res[(k+1)\%\ chnum]] - \\ p[res[i]]) > fabs(Cross(p[res[j]] - p[res[i]], \\ \end{array} 
          p[res[k]] - p[res[i]]))) k = (k + 1) % chnum;
     tmp = fabs (Cross (
     p[res[i]] - p[res[i]], p[res[k]] - p[res[i]]));
if (tmp > area) area = tmp;
     while (fabs(Cross(p[
         p[res[k]] - p[res[i]]))) j = (j + 1) % chnum;
     tmp = fabs (Cross (
         p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
     if (tmp > area) area = tmp;
  return area / 2;
8.9 Point in Polygon [0a9a66]
int pip(vector<P> ps, P p) {
  for (int i = 0; i < ps.size(); ++i) {
     int a = i, b = (i + 1) \% ps.size();
    L l(ps[a], ps[b]);
    P q = l.project(p);
     if \ ((p - q).abs() < eps \ \& \ l.inside(q)) \ return \ 1;
     if (same(ps[
         a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
     if (ps[a].y > ps[b].y) swap(a, b);
if (ps[a].y \le p.y && p.y <
          ps[b].y \& p.x \le ps[a].x + (ps[b].x - ps[a].x
) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) ++c;
```

8.10 Circle [466c44]

 $d = \max(0., d);$

double i = (-B - sqrt(d)) / (2 * A);

return (c & 1) * 2;

```
struct C {
 P c;
  double r;
 C(P \ c = P(0, 0), \ double \ r = 0) : c(c), r(r) \ \{\}
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r,
       d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
    double o = acos
        ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
    P i = (b.c - a.c).unit();
    p.push\_back(a.c + i.rot(o) * a.r);
    p.push\_back(a.c + i.rot(-o) * a.r);
  return p;
double IntersectArea (C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d \ge a.r + b.r - eps) return 0;
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos
      ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
  double q = acos
 // remove second
     level if to get points for line (defalut: segment)
vector <P> Circle Cross Line (Pa, Pb, Po, double r) {
 double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y); double C = sq(a.x - o.x)
      ) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
  if (d \ge -eps)
```

```
\begin{array}{l} \mbox{double } j = (-B + sqrt(d)) \ / \ (2 \ ^*A); \\ \mbox{if } (i \ - 1.0 <= eps \ \& \ i >= \end{array}
         -eps) t.emplace\_back(a.x + i * x, a.y + i * y);
    if (j - 1.0 \le eps \&\& j > =
         -eps) t.emplace_back(a.x + j * x, a.y + j * y);
  }
  return t;
// calc area
      intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(Pa, Pb, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
    return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
  if (inb) return
  SectorArea(p[0], a, r) + abs(p[0] \hat{} b) / 2;
if (p.size() = 2u) return SectorArea(a, p[0], r)
+ SectorArea(p[1], b, r) + abs(p[0] \hat{} p[1])
                                                     p[1]) / 2;
  else return SectorArea(a, b, r);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
    int j = (i + 1) \% 3;
    double o = atan2
    ans \ +\!\!= \ AreaOfCircleTriangle
         (ps[i], ps[j], r) * (o>= 0 ? 1 : -1);
  return abs(ans);
```

8.11 Tangent of Circles and Points to Circle [19eb58]

```
vector <\!\!L\!\!> tangent (C~a,~C~b)~\{
#define Pij \
 P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
  z.emplace\_back(a.c + i, a.c + i + j);
#define deo(I,J)\
  double d = (a)
      .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
 if (a.r < b.r) swap(a, b);
  vector<L> z;
  if ((a.c - b.c).abs() + b.r < a.r) return z;
  else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
  else
    deo(-,+);
    if (same(d, a.r + b.r)) { Pij; } else if (d > a.r + b.r) { deo(+,-); }
  return z;
}
vector<L> tangent(C c, P p) {
  vector < L > z;
  double d = (p - c.c).abs();
  if (same(d, c.r)) {
   P i = (p - c.c).rot(pi / 2);
    z.emplace\_back(p, p + i);
  } else if (\overline{d} > c.r) {
    \frac{\text{double o} = a\cos(c.r / d);}
    P i = (p - c.c).unit
        (), j = i.rot(o) * c.r, k = i.rot(-o) * c.r;
   z.emplace_back(c.c + j, p);
z.emplace_back(c.c + k, p);
  return z;
```

8.12 Area of Union of Circles [490636]

```
vector<pair<double, double>>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
  vector<pair<double, double>>> res;
  if (same(a.r + b.r, d));
  else if (d <= abs(a.r - b.r) + eps) {</pre>
```

```
if (a.r < b.r) res.emplace_back(0, 2 * pi);
    else if (d < abs(a.r + b.r) - eps) {
    If (z < 0) z = 2 p_1,

double l = z - 0, r = z + 0;

if (l < 0) l += 2 * p_1;

if (r > 2 * p_1) r == 2 * p_1;

if (l > r) res.emplace_back

(l, 2 * p_1), res.emplace_back(0, r);
    else res.emplace_back(l, r);
  return res;
double CircleUnionArea
    int n = c.size();
  double a = 0, w;
for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>>> s = {{2 * pi, 9}}, z; for (int j = 0; j < n; ++j) if (i != j) {
   z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    for (auto &e : s) {
      if (e.first > w) a += F(e.first) - F(w);
      w = max(w, e.second);
  return a * 0.5;
```

8.13 Minimun Distance of 2 Polygons [7eb8bb]

8.14 2D Convex Hull [65eaab]

```
bool operator < (const P &a, const P &b) {
 return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator>(const P &a, const P &b) {
  return same(a.x, b.x) ? a.y > b.y : a.x > b.x;
#define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return
  same(a.x, b.x) ? a.y < b.y : a.x < b.x; });
for (int i = 0; i < ps. size(); ++i) {
     while (p.size() >= 2 \&\& crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() -
2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push\_back(ps[i]);
```

```
p.pop_back();
                                                                      while (l + 1 < r) {
  return p;
                                                                        int m = (1 + r) >> 1;
int sgn(double
                                                                        else r = m;
     x) { return same(x, 0) ? 0 : x > 0 ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
                                                                    vector <P> getIS(P a, P b) {
                                                                      if (X > Y) swap(X, Y)
struct CH {
  int n;
  vector < P > p, u, d;
                                                                      return { };
  CH() \{ \}
  CH(vector < P > ps) : p(ps) {
    n = ps.size();
    rotate (p. begin
         (), min_element(p.begin(), p.end()), p.end());
    auto t = max_{element}(p.begin(), p.end());
                                                                      if (l == r) return;
    d = vector < P > (p.begin(), next(t));
    u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find (vector <P> &v, P d) {
    int l = 0, r = v.size();
while (l + 5 < r) {
int L = (l * 2 + r) / 3, R = (l + r * 2) / 3;
if (v[L] * d > v[R] * d) r = R;
                                                                        int m = (l + r) \gg 1;
                                                                        if (sgn
                                                                        else r = m;
       else l = L;
    int x = 1;
    for (int i = l +
                                                                    bool contain (P p) {
          return x;
                                                                      auto it
  int findFarest(P v) {
                                                                      if (it->x = p.x) {
    if (v.y > 0 \mid | v.y = 0 & v.x > 0) return
          ((int)d.size() - 1 + find(u, v)) \% p.size();
    return find(d, v);
                                                                      it = lower bound
  P get(int 1, int r, Pa, Pb) {
                                                                      if (it->x = p.x) {
    int s = sgn(crx(a, b, p[l \% n]));
    while (1 + 1 < r) {
       int \dot{m} = (l + r) \gg 1;
                                                                      return 1;
        \begin{tabular}{l} if & (sgn(crx(a, b, p[m \% n])) == s) & l = m; \\ \end{tabular} 
       else r = m;
    return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
                                                                      a = b = 0;
  }
                                                                      int i
  vector<P> getLineIntersect(P a, P b) {
    int X = findFarest((b - a).rot(pi / 2));
int Y = findFarest((a - b).rot(pi / 2));
                                                                      bs(0, i, p, a, b);
                                                                      bs(i, d.size(), p, a, b);
    if (X > Y) swap(X, Y);
                                                                      i = lower_bound(
    if (sgn
          \begin{array}{l} (crx(a,\ b,\ p[X])) \ * \ sgn(crx(a,\ b,\ p[Y])) < 0) \\ return \ \{get(X,\ Y,\ a,\ b),\ get(Y,\ X+n,\ a,\ b)\}; \end{array} 
                                                                      bs((int
    return {}; // tangent case falls here
                                                                      bs((int)d.size()
  void update_tangent(P q, int i, int &a, int &b) {
                                                                      return 1;
    if (sgn(crx(q, p[a], p[i])) > 0) a = i;
    if (sgn(crx(q, p[b], p[i])) < 0) b = i;
                                                                    8.15
  void bs(int l, int r, P q, int &a, int &b) {
    if (l == r) return
                                                                   double
    update_tangent(q, 1 % n, a, b)
    int s = sgn(crx(q, p[1 \% n], p[(1 + 1) \% n]));
while (1 + 1 < r) {
       int m = (l + r) >> 1;
                                                                   struct convex3D {
       if (sgn(crx
           (q,\ p[m\ \%\ n]\ ,\ p[(m+1)\ \%\ n])) == s)\ l = m;
                                                                      struct T{
                                                                        int a,b,c;
                                                                        bool res;
    update\_tangent(q,\ r\ \%\ n,\ a,\ b);
                                                                        T(){}
                                                                        T(int a, int
  int x = 1;
  for (int i = 1)
        + 1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
                                                                      int n,m;
  return x;
                                                                      P p[maxn];
                                                                      T f [maxn* 8];
int findFarest(P v) {
                                                                      int id [maxn] [maxn]:
  if (v.y > 0 | | v.y = 0 \& v.x > 0) return
                                                                      bool on (T &t, P &q) {
        ((int)d.size() - 1 + find(u, v)) \% p.size();
                                                                        return (
  return find (d, v);
P get(int 1, int r, Pa, Pb) {
                                                                      void meow(int q,int a,int b){
```

```
int s = sgn(crx(a, b, p[l \% n]));
    if (\operatorname{sgn}(\operatorname{crx}(a, b, p[m \% n])) == s) l = m;
  return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
  int X = findFarest((b - a).spin(pi / 2));
   int Y = findFarest((a - b).spin(pi / 2)); 
  if \ (sgn(crx(a,\ b,\ p[X]))\ *\ sgn(crx(a,\ b,\ p[Y])) <
      0) return \{ get(X, Y, a, b), get(Y, X + n, a, b) \};
void update_tangent(P q, int i, int &a, int &b) {
  \begin{array}{ll} \mbox{if } \left( sgn\left( crx\left( q,\ p\left[ a\right] ,\ p\left[ i\right] \right) \right) > 0 \right) \ a = i \,; \end{array}
  if (sgn(crx(q, p[b], p[i])) < 0) b = i;
void bs(int l, int r, Pq, int &a, int &b) {
  update_tangent(q, 1 % n, a, b);
  int s = sgn(crx(q, p[1 % n], p[(1 + 1) % n])); while (1 + 1 < r) {
         (crx(q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
  update\_tangent(q, r \% n, a, b);
  if (p.x < d[0].x \mid | p.x > d.back().x) return 0;
      = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
    if (it->y > p.y) return 0;
    else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
       (u.begin(), u.end(), P(p.x, 1e12), greater < P > ());
    if (it->y < p.y) return 0;
  } else if (crx(*prev(it), *it, p) < -eps) return 0;
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
      = lower\_bound(d.begin(), d.end(), p) - d.begin();
      u.begin(), u.end(), p, greater < P > ()) - u.begin();
       d.size() - 1, (int)d.size() - 1 + i, p, a, b);
        -1 + i, (int)d.size() - 1 + u.size(), p, a, b);
        3D Convex Hull [29e4a9]
```

```
absvol(const P a, const P b, const P c, const P d) {
return abs(((b-a)^(c-a))^*(d-a))/6;
static const int maxn=1010;
       b, int c, bool res=1: a(a), b(b), c(c), res(res){}
      p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
```

if (norm2(cent - p[k]) <= r) continue;</pre>

```
int g=id[a][b];
                                                                                          cent = center(p[i], p[j], p[k]);
     if(f[g].res){
                                                                                          r = norm2(p[k] - cent);
        if(on(f[g],p[q]))dfs(q,g);
                                                                                    }
           id [q][b]=id[a][q]=id[b][a]=m;
           f[m++]=T(b,a,q,1);
                                                                                 return circle(cent, sqrt(r));
     }
                                                                               8.17 Closest Pair [f6de57]
   void dfs(int p,int i){
                                                                              double closest_pair(int l, int r) {
     f[i].res=0;
                                                                                 // p should be sorted
      \begin{array}{l} \text{meow}(p,f[\,i\,\,].\,b,f[\,i\,\,].\,a)\,;\\ \text{meow}(p,f[\,i\,\,].\,c\,,f[\,i\,\,].\,b)\,;\\ \text{meow}(p,f[\,i\,\,].\,a\,,f[\,i\,\,].\,c)\,;\\ \end{array} 
                                                                                         increasingly according to the x-coordinates.
                                                                                      (1 = r) return 1e9;
                                                                                 if (r - l = 1) return dist(p[l], p[r]);
                                                                                 int m = (l + r) >> 1;
   void operator()(){
                                                                                 double d =
     if (n<4)return;
                                                                                        min(closest\_pair(l, m), closest\_pair(m + 1, r));
     if ([&](){
                                                                                  vector<int> vec;
           for (int i=1; i< n; ++i) if (abs
                                                                                  for (int i = m; i >= 1 &&
                (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
                                                                                 \begin{array}{l} fabs\,(p\,[m].\,x\,-\,p\,[\,i\,].\,x)\,<\,d\,;\,\,\,\text{--}\,i\,)\,\,\,vec\,.\,push\_back\,(\,i\,)\,;\\ for\,\,\,(\,int\,\,\,i\,=\,m\,+\,\,1\,;\,\,i\,<=\,r\,\,\&\&\\ fabs\,(p\,[m].\,x\,-\,p\,[\,i\,].\,x)\,<\,d\,;\,\,+\!\!+\!\!i\,)\,\,\,vec\,.\,push\_back\,(\,i\,)\,; \end{array}
           return 1;
}() || [&](){
           for (int i=2; i< n; ++i) if (abs((p[0]-p[i])
                                                                                 sort(vec.begin(), vec.end()
    , [&](int a, int b) { return p[a].y < p[b].y; });</pre>
                 (p[1]-p[i]) > eps) return swap(p[2],p[i]),0;
           return 1
                                                                                  for (int i = 0; i < vec.size(); ++i) {
           }() | | [&](){
                                                                                    for (int j = i + 1; j < vec.size()
    && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {
    d = min(d, dist(p[vec[i]], p[vec[j]]));
           for (int i
                 =3;i < n;++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                 *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
           return 1;
                                                                                 }
           }())return;
                                                                                 return d;
     for (int i=0; i<4;++i){
                                                                              }
        T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
        if(on(t,p[i]))swap(t.b,t.c);
                                                                               9
                                                                                     Else
        id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
        f[m++]=t;
                                                                               9.1 Cyclic Ternary Search* [9017cc]
     for (int i=4; i< n; ++i) for
                                                                                  bool pred(int a, int b);
           (int j=0; j < m++j) if (f[j]. res & on(f[j], p[i])){
                                                                               f(0) \sim f(n-1) is a cyclic-shift U-function
        dfs\left( \,i\,\,,\,j\,\right) ;
                                                                               return idx s.t. pred(x, idx) is false for all x*/
        break;
                                                                               int cyc_tsearch(int n, auto pred) {
                                                                                 if (n = 1) return 0;
     int mm=m; m=0;
                                                                                 int l = 0, r = n; bool rv = pred(1, 0);
     for(int i=0;i<mm++i)if(f[i].res)f[m++]=f[i];
                                                                                 while (r - 1 > 1) {
    int m = (1 + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
  else l = m;
                                                                                 return pred(1, r % n) ? 1 : r % n;
           (p[f[i].a], p[f[i].b], p[f[i].c], p[f[j].c])>eps);
                                                                              9.2
                                                                                      Mo's
                                                                                                    Algorithm(With
                                                                                                                                    modification)
   int faces(){
     int r=0;
                                                                                       [f05c5b]
     for (int i=0; i < m++i) {
        int iden=1;
                                                                              Mo's Algorithm With modification
        \label{eq:for_int} \begin{array}{ll} \text{for} \, (\, \text{int} \  \, j \! = \! 0; j \! < \! i; \! + \! + \! j \,) \, \text{if} \, (\, \text{same} \, (\, i \, , \, j \,) \,) \, \text{iden} \! = \! 0; \end{array}
                                                                              Block: N^{2/3}, Complexity: N^{5/3}
        r + = iden;
                                                                               struct Query {
     return r;
                                                                                 int L, R, LBid, RBid, T;
Query(int l, int r, int t):
} tb;
                                                                                    L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
8.16 Minimum Enclosing Circle [fc0e72]
                                                                                 bool operator < (const Query &q) const {
  if (LBid != q.LBid) return LBid < q.LBid;
pt center(const pt &a, const pt &b, const pt &c) {
                                                                                    \quad \text{if} \quad (RBid \ != \ q.RBid) \quad \textbf{return} \quad RBid < \ q.RBid;
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
                                                                                    return T < b.T;
   double d = p0 \hat{p}1;
                                                                               };
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
                                                                               void solve(vector<Query> query) {
                                                                                 sort(ALL(query));
  return pt(x, y);
                                                                                 int L=0, R=0, T=-1;
                                                                                 for (auto q : query)
                                                                                    while (T < q.T) addTime(L, R, ++T); // TODO while (T > q.T) subTime(L, R, T--); // TODO
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
                                                                                    while (R < q.R) add(arr[++R]); // TODO
while (L > q.L) add(arr[--L]); // TODO
while (R > q.R) sub(arr[R--]); // TODO
   double r = 0.0;
   pt cent;
   for (int i = 0; i < p.size(); ++i) {
                                                                                    while (L < q.L) sub(arr[L++]); // TODO
     if (norm2(cent - p[i]) <= r) continue;
                                                                                    // answer query
     cent = p[i];
     r = 0.0:
                                                                              }
     for (int j = 0; j < i; +++j) {
        if (norm2(cent - p[j]) \le r) continue;
                                                                                      Mo's Algorithm On Tree [8331c2]
        cent = (p[i] + p[j])
        r = norm2(p[j] - cent);
        for (int k = 0; k < j; ++k) {
                                                                              Mo's Algorithm On Tree
```

Preprocess:

```
1) LCA
 2) \hspace{0.1in} \mathrm{dfs} \hspace{0.1in} \mathrm{with} \hspace{0.1in} \mathrm{in} \hspace{0.1in} [\hspace{0.1in} \mathrm{u}\hspace{0.1in}] \hspace{0.1in} = \hspace{0.1in} \mathrm{dft} \hspace{-0.1in} +\hspace{-0.1in} +, \hspace{0.1in} \mathrm{out} \hspace{0.1in} [\hspace{0.1in} \mathrm{u}\hspace{0.1in}] \hspace{0.1in} = \hspace{0.1in} \mathrm{dft} \hspace{-0.1in} +\hspace{-0.1in} +
 3) \operatorname{ord}[\operatorname{in}[\mathbf{u}]] = \operatorname{ord}[\operatorname{out}[\mathbf{u}]] = \mathbf{u}
     bitset ⟨MAXN⟩ inset
struct Query {
int L, R, LBid, lca;
     Query(int u, int v) {
         int c = LCA(u, v);
        if (c = u \mid c = v)
             q.\,lca \,=\, -1\,, \,\, q.\,L \,=\, out\,[\,c\,\, \widehat{}\,\, u\,\, \widehat{}\,\, v\,]\,\,, \,\, q.\,R \,=\, out\,[\,c\,\,]\,; 
         else if (out[u] < in[v])
            q.lca = c, q.L = out[u], q.R = in[v];
        \begin{array}{l} q.\,lca \, = \, c \, , \ q.L \, = \, out \, [v] \, , \ q.R \, = \, in \, [u] \, ; \\ q.\,Lid \, = \, q.L \, \, / \  \, blk \, ; \end{array}
     bool operator < (const Query &q) const {
         if (LBid != q.LBid) return LBid < q.LBid;
        return R < q.R;
    }
 f,
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
    else add(arr[x]); // TODO
        inset[x] = \sim inset[x];
 void solve(vector<Query> query) {
     sort (ALL(query));
     int \dot{L} = \dot{0}, R = 0;
     for (auto q : query) {
        while (R < q.R) flip (ord[++R]);
         while (L > q.L) flip (ord[--L]);
        while (R > q.R) flip (ord[R--]); while (L < q.L) flip (ord[L++]);
        if (~q.lca) add(arr[q.lca]);
// answer query
         if (~q.lca) sub(arr[q.lca]);
    }
| }
            Additional Mo's Algorithm Trick
  • Mo's Algorithm With Addition Only
       - \ \ Sort querys same as the normal Mo's algorithm.
       - For each query [l,r]:
       - If l/blk = r/blk, brute-force.
       - If l/blk \neq curL/blk, initialize curL := (l/blk + 1) \cdot blk, curR :=
           curL-1
       - If r > curR, increase curR
       - decrease curL to fit l, and then undo after answering
  • Mo's Algorithm With Offline Second Time
       - \ \operatorname{Require} : \operatorname{Changing} \operatorname{answer} \equiv \operatorname{adding} f([l,r],r+1).
       - Require: f([l,r],r+1) = f([1,r],r+1) - f([1,l),r+1).
- Part1: Answer all f([1,r],r+1) first.
           Part2: Store curR \to R for curL (reduce the space to O(N)), and then
           answer them \, by \, the \, second \, off line \, algorithm.
                   You must do the above symmetrically for the left boundaries.
```

9.5 Hilbert Curve [1274a3]

```
ll hilbert(int n, int x, int y) {
  11 \text{ res} = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y \& s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry = 0) {
       if (rx = 1) x = s - 1 - x, y = s - 1 - y;
      swap(x, y);
    }
  }
  return res;
 // n = 2^k
```

9.6 DynamicConvexTrick* [673ffd]

```
// only works for integer coordinates!! maintain max struct Line {
  mutable ll a, b, p;
  bool operator
      <(const Line &rhs) const { return a < rhs.a; }
  bool operator < (ll x) const { return p < x; }
struct DynamicHull : multiset<Line, less<>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
        lì b) { return a / b - ((a \hat{b}) < 0 \& a \% b); }
  bool isect(iterator x, iterator y) {
```

```
if (y = end()) \{ x->p = kInf; return 0; \}
     if (x
          ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
     else x->p = Div(y->b - x->b, x->a - y->a);
     return x->p>= y->p;
  void addline(ll a, ll b) {
       \  \, \text{auto} \ z \, = \, insert \, (\, \{a \, , \ b \, , \ \dot{0} \, \}) \, , \ y \, = \, z + +, \ x \, = \, y \, ; 
     while (isect(y, z)) z = erase(z);
     if (x != begin
          () && isect(--x, y)) isect(x, y = erase(y));
     while ((y = x) != begin
          () && (-x)-p >= y-p) isect(x, erase(y));
  ll query(ll x) {
     auto l = *lower_bound(x);
     return l.a * x + l.b;
};
      All LCS* [78a378]
9.7
\begin{tabular}{ll} void & all\_lcs(string s, string t) & $//$ 0-base \\ \end{tabular}
  vector < int > h(SZ(t));
  iota(ALL(h), 0);
  for (int a = 0; a < SZ(s); ++a) {
```

```
int v = -1;
for (int c = 0; c < SZ(t); ++c)
   if'(s[a] = t[c] | | h[c] < v)
     swap(h[c], v)
// LCS(s[0, a], t[b, c]) =
// c - b + 1 - sum([h[i] >= b] | i <= c)
// h[i] might become -1 !!
```

9.8 AdaptiveSimpson* [4074b3]

}

```
template<typename Func, typename d = double>
struct Simpson {
   \begin{array}{ll} \textbf{using} & \textbf{pdd} = \textbf{pair} < \!\! d\,, \ d \!\! >; \end{array}
   Func f;
  pdd mix(pdd 1, pdd r, optional<d> fm = {}) {
    d h = (r.X - 1.X) / 2, v = fm.value_or(f(1.X + h));
    return {v, h / 3 * (1.Y + 4 * v + r.Y)};
   d eval(pdd 1, pdd r, d fm, d eps) {
  pdd m((1.X + r.X) / 2, fm);
  d s = mix(1, r, fm).second;
      auto [flm, sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
      d \ delta = sl + sr - s;
      if (abs(delta
             ) <= 15 * eps) return sl + sr + delta / 15;
      return eval(1, m, flm, eps / 2) +
         eval(m, r, fmr, eps / 2);
   d eval(d l, d r, d eps) {
      return eval
             (\{1, f(1)\}, \{r, f(r)\}, f((1+r) / 2), eps);
  d eval2(d l, d r, d eps, int k = 997) {
d h = (r - l) / k, s = 0;
for (int i = 0; i < k; ++i, l += h)
         s \leftarrow eval(l, l + h, eps / k);
      return s;
};
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
```

9.9 Simulated Annealing [de78c6]

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times for (int it = 1; it <= 1000000; ++it) {
          answer, nw: current value, rnd(): mt19937 rnd()
     if (\exp(-(nw - ans))
          ) / factor) >= (double)(rnd() % base) / base)
         ans = nw:
     factor *= 0.99995;
```

9.10 Tree Hash* [34aae5]

```
ull seed;
ull shift (ull x) {
```

```
National Yang Ming Chiao Tung University FubukiMyWife
     \hat{} = x << 13;
  x = x >> 7;
  x = x << 17;
  return x;
ull dfs(int u, int f) {
  ull sum = seed;
  for (int i : G[u])
     if (i != f)
       sum += shift(dfs(i, u));
  return sum;
9.11 Binary Search On Fraction [765c5a]
struct Q {
  ll p, q
  Q go(Qb, ll d) \{ return \{p + b.p*d, q + b.q*d\}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo\{0, 1\}, hi\{1, 0\};
  if (pred(lo)) return lo;
  assert(pred(hi));
bool dir = 1, L = 1, H = 1;
for (; L || H; dir = !dir) {
     ll len = 0, step = 1;
for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
        \begin{array}{l} \text{if } (Q \text{ mid} = \text{hi.go}(\text{lo}, \text{len} + \text{step}); \\ \text{mid.p} > N \mid \mid \text{mid.q} > N \mid \mid \text{dir } \hat{} \text{pred}(\text{mid})) \end{array}
          t++;
        else len += step;
     swap(lo, hi = hi.go(lo, len));
     (dir ? L : H) = !!len;
  return dir ? hi : lo;
9.12 Min Plus Convolution* [09b5c3]
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution
     (vector < int > &a, vector < int > &b) {
  int n = SZ(a), m = SZ(b);
   \begin{array}{l} \text{vector} < \text{int} > c \left( n + m - 1, \text{ INF} \right); \\ \text{auto} \ dc = [\&] \left( \text{auto } Y, \text{ int } 1, \text{ int } r, \text{ int } \text{jl}, \text{ int } \text{jr} \right) \end{array} \} 
     if (l > r) return;
     int mid = (1 + r) / 2, from = -1, &best = c[mid];
     for (int j = jl; j \leftarrow jr; ++j)
        if(int i = mid - j; i >= 0 && i < n)
          if (best > a[i] + b[j])

best = a[i] + b[j], from = j;
     Y(Y, l,
          mid - 1, jl, from), Y(Y, mid + 1, r, from, jr);
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
9.13 Bitset LCS [330ab1]
cin >> n >> m;
for (int i = 1, x; i \le n; ++i)
  cin \gg x, p[x].set(i);
for (int i = 1, x; i \le m; i++) {
cin >> x, (g = f) |= p[x];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) = g) \&= g;
cout << f.count() << '\n';
9.14 N Queens Problem [d1fccc]
     (vector<int> &ret, int n) { // no sol when n=2,3
  if (n % 6 == 2) {
for (int i = 2; i <= n; i += 2) ret.pb(i);
     ret.pb(3); ret.pb(1);
     for (int i = 7; i \le n; i += 2) ret.pb(i);
```

ret.pb(5);

ret.pb(2);

else {

ret.pb(1); ret.pb(3);

} else if (n % 6 == 3) {
 for (int i = 4; i <= n; i += 2) ret.pb(i);</pre>

for (int i = 5; $i \le n$; i += 2) ret.pb(i);

for (int i = 2; $i \le n$; i += 2) ret.pb(i);

```
10 Python
10.1 Misc
# decimal
from decimal import *
setcontext (Context (prec
=MAX_PREC, Emax=MAX_EMAX, rounding=ROUND_FLOOR))
print(Decimal(input()) * Decimal(input()))
from fractions import Fraction
Fraction
     ("3.14159").limit_denominator(10).numerator # 22
# N*M
map(int, input().split())
arr2d = [[
     list(map(int, input().split()))] for i in range(N)]
# a^b%M
pow(a, b, M)
# random
from random import *
randrange(L, R, step) # [L,R) L+k*step
randint(L, R) # int from [L,R]
choice(list) # pick 1 item from list
choices(list, k) # pick k item
shuffle (list)
Uniform(L, R) # float from [L,R]
# print
print(f"num: {num}, pi: {pi:.2f}, str: {text}")
print(1, 2, 3, 4, sep=" | ", end=" <-- END\n")</pre>
# file IO
r = open("filename.in")
a = r.read() # read whole content into one string
w = open("filename.out", "w")
w. write("123\n")
# IO redirection
import sys
sys.stdin = open("filename.in")
sys.stdout = open("filename.out", "w")
print("123", file=sys.stderr)
input\_data = sys.stdin.buffer.read().split()
n = int(input_data[0])
arr = list(map(int, input_data[1 : n + 1]))
sys.set\_int\_max\_str\_digits(5000000)
10.2 Stress testing
i = 0
while true
do
    i=\$((\$i+1))
    python3 gen.py > input.txt
python3 brute.py < input.txt > brute.txt
     ./sol < input.txt > output.txt
     diff output.txt brute.txt
    if [ $? -ne 0 ]
    then
         break
     fi
done
```

for (int i = 1; $i \le n$; i += 2) ret.pb(i);