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#### 1.1 vimrc

```
"This file should be placed at ~/.vimrc" se nu ai hls et ru ic is sc cul se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a syntax on hi cursorline cterm=none ctermbg=89 set bg=dark inoremap {<CR> {<CR>}<Esc>ko<tab> "Select region and then type :Hash to hash your selection." "Useful for verifying that there aren't mistypes." ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \| md5sum \| cut -c-6
```

# 1.2 readchar [0754b0] inline char readchar() {

static const size\_t bufsize = 65536;

```
static char buf[bufsize];
static char *p = buf, *end = buf;
  if (p == end) end = buf +
       fread_unlocked(buf, 1, bufsize, stdin), p = buf;
  return *p++;
}
1.3 Black Magic [d566f1]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
     tree<int, null_type, std::less<int>, rb_tree_tag
      tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
      (0) = 22; assert(*s.find_by_order(1) = 71);
  assert(s.order_of_key
      (22) = 0; assert (s.order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by_order
      (0) = 71; assert (s.order_of_key(71) = 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope < char > r[2];
  r[1] = r[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0;
1.4 Pragma Optimization [7b330a]
```

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse, sse2, sse3, sse3, sse4")
#pragma GCC target("popent, abm, nmmx, avx, arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

#### 1.5 Bitset [282252]

```
#include < bits / stdc++.h>
using namespace std;

int main () {
    bitset <4> bit;
    bit.all(); // all bit is true, ret tru;
    bit.any(); // any bit is true, ret true
    bit.none(); // all bit is false, ret true
    bit.count();
    bit.to_string('0', '1'); // with parmeter
    bit.reset(); // set all to true
    bit.set(); // set all to false
    std::bitset <8> b3{0}, b4{42};
    std::hash<std::bitset <8> hash_fn8(b4);
}
```

# 2 Graph

### $2.1\quad BCC\ Vertex*\ [ed8308]$

```
struct BCC { // 0-base
  int n, dft, nbcc;
  vector<int> low, dfn, bln, stk, is_ap, cir;
  vector<vector<int>>> G, bcc, nG;
  void make_bcc(int u) {
    bcc.emplace_back(1, u);
    for (; stk.back() != u; stk.pop_back())
       bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
    stk.pop_back(), bln[u] = nbcc++;
```

for (; stk.back() != u; stk.pop\_back())

```
bln [stk
  void dfs(int u, int f) {
                                                                                                      .back()] = nscc, instack[stk.back()] = 0;
                                                                                            instack[u] = 0, bln[u] = nscc++, stk.pop_back();
     int child = 0;
     low[u] = dfn[u] = ++dft, stk.pb(u);
     for (int v : G[u])
        if (!dfn[v]) {
                                                                                      \begin{array}{l} dfs\left(v,\;u\right),\; +\!\!\!\!+\!\!\!\!+\!\!\!\!\!+ child\;;\\ low\left[u\right]\; =\; min\left(low\left[u\right]\;,\; low\left[v\right]\right)\;; \end{array}
                                                                                      void add_edge(int u, int v) {
           if (dfn[u] \le low[v]) {
 is\_ap[u] = 1, bln[u] = nbcc;
                                                                                         G[u].pb(v);
                                                                                      make\_bcc(v), bcc.back().pb(u);
                                                                                            if (!dfn[i]) dfs(i);
        else if (dfn[v] < dfn[u] & v != f)
     low[u] = min(low[u], dfn[v]);
if (f = -1 & child < 2) is ap[u] = 0;
                                                                                   }; // scc_id(i): bln[i]
     if (f = -1 & child = 0) make_bcc(u);
                                                                                   2.4 2SAT* [e839e5]
   \begin{array}{lll} \dot{B}CC(int \ \underline{\ \ } n): \ n(\underline{\ \ \ } n) \ , \ dft \, () \ , \\ nbcc \, () \ , \ low \, (n) \ , \ dfn \, (n) \ , \ bln \, (n) \ , \ is\_ap \, (n) \ , \ G(n) \ \ \{\} \\ void \ add\_edge \, (int \ u, \ int \ v) \ \ \{ \end{array} 
                                                                                   struct SAT { // 0-base
                                                                                      int n:
                                                                                      vector<bool> istrue;
     G[u].pb(v), G[v].pb(u);
                                                                                      SCC scc;
                                                                                      SAT(int \underline{n}): n(\underline{n}), istrue(n+n), scc(n+n)  {}
  void solve() {
  for (int i = 0; i < n; ++i)</pre>
                                                                                      int rv(int a) {
                                                                                         return a > = n ? a - n : a + n;
        if (!dfn[i]) dfs(i, -1);
                                                                                      void add_clause(int a, int b) {
  void block_cut_tree() {
                                                                                         scc.add\_edge(rv(a), b), scc.add\_edge(rv(b), a);
     cir.resize(nbcc);
     for (int i = 0; i < n; ++i)
        if (is_ap[i])
bln[i] = nbcc++;
                                                                                      bool solve()
                                                                                         scc.solve();
                                                                                         for (int i = 0; i < n; ++i) {
     cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
                                                                                            if(scc.bln[i] = scc.bln[i+n]) return false;
                                                                                            istrue[i] = scc.bln[i] < scc.bln[i + n];
        for (int j : bcc[i])
                                                                                            istrue[i + n] = !istrue[i];
           if (is_ap[j])
  nG[i].pb(bln[j]), nG[bln[j]].pb(i);
} // up to 2 * n - 2 nodes!! bln[i] for id
                                                                                         return true;
                                                                                      }
                                                                                   };
2.2 Bridge* [Occada]
                                                                                           MinimumMeanCycle* [4be648]
struct ECC { // 0-base
  int n, dft, ecnt, necc;
vector<int> low, dfn, bln, is_bridge, stk;
                                                                                   ll road[N][N]; // input here
                                                                                   struct MinimumMeanCycle {
  {\tt vector}{<\hspace{-.05cm}{\rm vic}}{\times}{\rm G};
                                                                                      11 \, dp[N + 5][N], n;
                                                                                       pll solve() {
  void dfs(int u, int f) {
                                                                                         11 a = -1, b = -1, L = n + 1;
for (int i = 2; i \le L; ++i)
     dfn\left[u\right] \,=\, low\left[u\right] \,=\, +\!\!+\!\! dft\;,\;\; stk\,.\,pb(u)\;;
     for (auto [v, e] : G[u])
if (!dfn[v])
                                                                                            for (int k = 0; k < n; ++k)
        \begin{array}{l} dfs\left(v,\,e\right),\;low\left[u\right] \,=\, min(low\left[u\right],\;low\left[v\right])\,;\\ else\;\;if\;\left(e\;!=\;f\right) \end{array}
                                                                                               for (int j = 0; j < n; ++j)
                                                                                                  dp[i][j] =
                                                                                          \min_{\substack{m \text{ in } (dp[i-1][k] + road[k][j], \ dp[i][j]); \\ \text{for } (int \ i=0; \ i< n; ++i) \ \{ } 
     \begin{array}{c} low\left[u\right] \stackrel{.}{=} min(low\left[u\right], \ dfn\left[v\right]);\\ if \ (low\left[u\right] \stackrel{.}{==} dfn\left[u\right]) \ \{ \end{array}
                                                                                            if (dp[L][i] >= INF) continue;
         if (f!=-1) is_bridge[f]=1;
                                                                                            for (; stk.back() != u; stk.pop_back())
bln[stk.back()] = necc;
        bln[u] = necc++, stk.pop\_back();
                                                                                                  ta \, = \, \dot{d}p \, [L] \, [\, \dot{i}\, ] \  \, - \, dp \, [\, \dot{j}\, ] \, [\, \dot{i}\, ] \, , \  \, tb \, = \, L \, \, - \, \, \dot{j} \, ;
                                                                                            if (ta == 0) continue;
  ÉCC(int _n): n(_n), dft()
         , \ \underline{ecnt}() \ , \ \underline{necc}() \ , \ low(\underline{n}) \ , \ dfn(\underline{n}) \ , \ bln(\underline{n}) \ , \ G(\underline{n}) \ \{\}
                                                                                            if (a = -1) | a * tb > ta * b) a = ta, b = tb;
  void add_edge(int u, int v)
                                                                                         if (a != -1) \{

11 g = gcd(a, b);
     G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
                                                                                            return pll(a / g, b / g);
  void solve() {
     is_bridge.resize(ecnt);
for (int i = 0; i < n; ++i)</pre>
                                                                                         return pll(-1LL, -1LL);
        if (!dfn[i]) dfs(i, -1);
                                                                                      void init(int _n) {
}; // ecc_id(i): bln[i]
                                                                                         for (int i = 0; i < n; ++i)
2.3 SCC* [22afe1]
                                                                                             \mbox{for (int } j \, = \, 0 \, ; \ j \, < \, n \, ; \ + + j \, ) \ dp [\, i \, + \, 2 \, ][\, j \, ] \, = \, INF \, ; 
struct SCC { // 0-base
                                                                                   };
  int n, dft, nscc;
  \label{eq:vector} vector <\!\!int\!\!> low\,, \ dfn\,, \ bln\,, \ instack\,, \ stk\,;
                                                                                   2.6 Virtual Tree* [80f7cb]
  vector < vector < int >>> G;
  void dfs(int u)
                                                                                   vector < int > vG[N];
     low[u] = dfn[u] = ++dft;
                                                                                    int top, st[N];
     instack[u] = 1, stk.pb(u);
     for (int v : G[u])
                                                                                   void insert(int u) {
        if^{(dfn[v])}
                                                                                      if (top = -1) return st[++top] = u, void();
        d\hat{f}s(v), low[u] = min(low[u], low[v]);
else if (instack[v] && dfn[v] < dfn[u])
                                                                                      int p = LCA(st[top], u);
                                                                                       if (p = st[top]) return st[++top] = u, void();
                                                                                      while (top >= 1 \&\& dep[st[top - 1]] >= dep[p])
          low[u] = min(low[u], dfn[v]);
     if (low[u] = dfn[u]) {
                                                                                      vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
```

dst[i][i] = vcst[i] = 0;

```
void chmin(int &x, int val) {
     vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
                                                                                   x = \min(x, val);
                                                                                 void add_edge(int ui, int vi, int wi) {
void reset(int u) {
                                                                                   chmin(dst[ui][vi], wi);
  for (int i : vG[u]) reset(i);
  vG[u].clear();
                                                                                void shortest_path() {
                                                                                   for (int k = 0; k < n; ++k)
                                                                                      for (int i = 0; i < n; ++i)
                                                                                        for (int j = 0; j < n; ++j)
void solve(vector<int> &v) {
                                                                                           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  top = -1
  sort (ALL(v),
     [\&](int a, int b) \{ return dfn[a] < dfn[b]; \});
                                                                                int solve(const vector<int>& ter) {
  for (int i : v) insert(i);
                                                                                   shortest_path();
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
                                                                                   int t = SZ(ter), full = (1 << t) - 1;
  // do something
                                                                                   for (int i = 0; i \leftarrow full; ++i)
                                                                                     fill_n (dp[i], n, INF);
  reset(v[0]);
                                                                                   \begin{array}{l} \operatorname{copy\_n}(\operatorname{vcst},\ \operatorname{n},\ \operatorname{dp}[\ 0\ ]);\\ \operatorname{for}\ (\operatorname{int}\ \operatorname{msk}=\ 1;\ \operatorname{msk}<=\ \operatorname{full};\ +\!\!+\!\!\operatorname{msk})\ \{ \end{array}
       Maximum Clique Dyn* [4a6b3d]
                                                                                      if (!(msk & (msk - 1))) {
                                                                                         \begin{array}{l} \text{int who} = \underline{\hspace{0.5cm}} \lg(msk); \\ \text{for (int } i = 0; i < n; +\!\!\!+\!\!\! i) \end{array} 
struct MaxClique { // fast when N \le 100 bitset A > G[N], cs[N]; int ans, sol[N], q, cur[N], d[N], n; void init (int _n) {
                                                                                           dp [msk
                                                                                                [i] = vcst[ter[who]] + dst[ter[who]][i];
     n = n;
                                                                                      for (int i = 0; i < n; ++i)
      \  \, \text{for} \  \, (\, \text{int} \  \, i \, = \, 0\,; \  \, i \, < \, n\,; \, +\!\!\!\!+\!\!\! i\,) \, \, G[\, i\, ]\,.\, reset\,(\,)\,; \\
                                                                                        for (int sub = (
                                                                                              msk - 1) & msk; sub; sub = (sub - 1) & msk)
  void add_edge(int u, int v) {
                                                                                           \begin{array}{c} chmin(dp[msk][i],\\ dp[sub][i] + dp[msk \ \widehat{\ } sub][i] - vcst[i]); \end{array}
    G[u][v] = G[v][u] = 1;
                                                                                      for (int i = 0; i < n; ++i) { tdst[i] = INF;
  \begin{tabular}{ll} void $$ pre\_dfs(vector<int> \&r, int 1, bitset< N> mask) $$ \{$ \end{tabular}
     if (1 < 4) {
                                                                                         for (int j = 0; j < n; ++j)
        for (int i : r) d[i] = (G[i] & mask).count();
                                                                                           chmin(tdst[i], dp[msk][j] + dst[j][i]);
        sort (ALL(r)
             , [\&](int x, int y) \{ return d[x] > d[y]; \});
                                                                                      copy_n(tdst, n, dp[msk]);
     vector < int > c(SZ(r));
                                                                                   return *min_element(dp[full], dp[full] + n);
     int 1 \text{ft} = \max(\text{ans} - q + 1, 1), \text{ rgt} = 1, \text{ tp} = 0;
     cs[1].reset(), cs[2].reset();
                                                                             ; // O(V 3^T + V^2 2^T)
     for (int p : r) {
                                                                              2.9 Dominator Tree* [e95beb]
        int k = 1;
        while ((cs[k] & G[p]).any()) ++k;
        if (k > rgt) cs[++rgt + 1].reset();
cs[k][p] = 1;
if (k < lft) r[tp++] = p;</pre>
                                                                              struct dominator_tree
                                                                                                              // 1-base
                                                                                vector<int> G[N], rG[N]; int n, pa[N], idn[N], id [N], Time; int semi[N], idom[N], best[N]; vector<int> tree[N]; // dominator_tree
     for (int k = lft; k \ll rgt; ++k)
                                                                                 void init(int _n) {
        for (int p = cs[k]._Find_first
                                                                                  n = _n;
for (int i = 1; i <= n; ++i)
rG[i]. clear(
          (); p < N; p = cs[k]._Find_next(p))
r[tp] = p, c[tp] = k, ++tp;
                                                                                     G[i].clear(), rG[i].clear();
     dfs(r, c, l + 1, mask);
                                                                                 void add_edge(int u, int v) {
  void dfs (vector<
                                                                                   G[u].pb(v), rG[v].pb(u);
        int>&r, vector<int>&c, int 1, bitset<N> mask) {
     while (!r.empty()) {
                                                                                void dfs(int u)
        int p = r.back();
                                                                                   id[dfn[u] = ++Time] = u;
        r.pop\_back(), mask[p] = 0;
                                                                                   for (auto v : G[u])
        if (q + c.back() <= ans) return;</pre>
                                                                                      if(!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
        \operatorname{cur}\left[q++\right] = p;
        vector<int> nr;
                                                                                int find(int y, int x) {
        for (int i : r) if (G[p][i]) nr.pb(i);
                                                                                   if (y \le x) return y;
        if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
                                                                                   int tmp = find(pa[y], x);
        else if (q > ans) ans = q, copy_n(cur, q, sol); c.pop_back(), --q;
                                                                                   if (semi[best[y]] > semi[best[pa[y]]])
                                                                                      best[y] = best[pa[y]];
     }
                                                                                   return pa[y] = tmp;
  int solve() {
                                                                                void tarjan(int root) {
     vector < int > r(n);
                                                                                   Time = 0;
     ans = q = 0, iota(ALL(r), 0);
                                                                                   for (int i = 1; i \le n; ++i) { dfn[i] = idom[i] = 0;
     pre\_dfs(r, 0, bitset<N>(string(n, '1')));
     return ans;
                                                                                      tree[i].clear();
                                                                                      best[i] = semi[i] = i;
};
                                                                                   dfs(root);
for (int i = Time; i > 1; --i) {
        Minimum Steiner Tree* [cbf811]
struct SteinerTree {
                                                                                      int u = id[i];
for (auto v : rG[u])
                            // 0-base
  int n, dst[N][N], dp[1 \ll T][N], tdst[N]; int vcst[N]; // the cost of vertexs
                                                                                        \inf (v = dfn[v]) \{
                                                                                           find(v, i);

semi[i] = min(semi[i], semi[best[v]]);
  void init(int _n) {
     for (int i = 0; i < n; ++i) {
        fill_n(dst[i], n, INF);
                                                                                      tree [semi[i]].pb(i);
```

for (auto v: tree[pa[i]]) {

find(v, pa[i]);idom[v] =

```
semi\left[\,best\left[\,v\,\right]\,\right] \; = \; pa\left[\,i\,\right] \;\;?\;\; pa\left[\,i\,\right] \;\; : \;\; best\left[\,v\,\right];
          tree [pa[i]]. clear();
      for (int i = 2; i \leftarrow Time; ++i) {
          if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
          tree [id [idom[i]]].pb(id[i]);
};
```

#### Minimum Clique Cover\* [5951ca] 2.10

```
\begin{array}{lll} & \textbf{struct} & Clique\_Cover \; \{ \; \; // \; 0\text{-base} \,, \; O(n2 \hat{}^{}n) \\ & \textbf{int} \; \; co[\; 1 << N] \,, \; n, \; E[N] \,; \\ & \textbf{int} \; \; dp[\; 1 << N] \,; \end{array}
    void init (int _n) {
    n = _n, fill_n (dp, 1 << n, 0);

        fill_n(E, n, 0), fill_n(co, 1 << n, 0);
    int solve() {
  for (int i = 0; i < n; ++i)
            co[1 << i] = E[i] | (1 << i);
        co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
        for (int i = 1; i < (1 << n); ++i) {
            or (int i -
int t = i & -i;
dn[i ^ t];
            \begin{array}{l} dp[\,i\,] \,=\, -dp[\,i\,\,\widehat{\ }\,\, t\,\,]\,;\\ co\,[\,i\,\,] \,=\, co\,[\,i\,\,\widehat{\ }\,\, t\,\,]\,\,\&\,\, co\,[\,t\,\,]\,; \end{array}
        for (int i = 0; i < (1 << n); ++i)

co[i] = (co[i] & i) == i;
         fwt(co, 1 << n, 1);
        for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)
    sum += (dp[i] *= co[i]);</pre>
             if (sum) return ans;
        return n;
    }
};
```

#### NumberofMaximalClique\* [c163d7]

```
struct BronKerbosch { // 1-base
   \begin{array}{l} int \ n, \ a\left[N\right], \ g\left[N\right]\left[\stackrel{.}{N}\right]^{'}, \\ int \ S, \ all \left[N\right]\left[\stackrel{.}{N}\right], \ some \left[N\right]\left[N\right], \ none \left[N\right]\left[N\right]; \end{array}
    void init(int _n) {
       n = \underline{n};
       for (int i = 1; i \le n; ++i)
           for (int j = 1; j \le n; ++j) g[i][j] = 0;
    void add_edge(int u, int v) {
       g[u][v] = g[v][u] = 1;
    void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
       if (sn = 0 \&\& nn = 0) + S;
       int u = some[d][0];
       for (int i = 0; i < sn; ++i) {
           int v = some[d][i];
          int v = some[d][1];
if (g[u][v]) continue;
int tsn = 0, tnn = 0;
copy_n(all[d], an, all[d + 1]);
all[d + 1][an] = v;
for (int j = 0; j < sn; ++j)
    if (g[v][some[d][j]])</pre>
           some [d + 1][tsn++] = some[d][j];
for (int j = 0; j < nn; ++j)
               if (g[v][none[d][j]])
                  none[d + 1][tnn++] = none[d][j];
           \begin{array}{l} dfs\,(d+1,\;an+1,\;tsn\,,\;tnn)\,;\\ some\,[d]\,[\,i\,]\,=\,0\,,\;none\,[d\,]\,[\,nn++]\,=\,v\,; \end{array}
    int solve() {
       iota(some[0], some[0] + n, 1);
       S = 0, dfs(0, 0, n, 0);
       return S;
};
```

#### 3 Data Structure

```
3.1 Discrete Trick
vector<int> val;
// build
sort (ALL
    (val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// \max idx <= :
upper_bound(ALL(val), x) - val.begin();
// \max idx < x
lower_bound(ALL(val), x) - val.begin();
3.2 BIT kth* [7d1b5f]
int bit [N + 1]; // N = 2 ^ k
int query_kth(int k) {
    int res = 0;
    for (int i = N >> 1; i >= 1; i >>= 1)
        if (bit [res + i] < k)
            k -= bit [res += i];
    return res + 1;
3.3 Interval Container* [78516e]
/* Add and
     remove intervals from a set of disjoint intervals.
 * Will merge the added interval with
     any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
\mathtt{set} \! < \! \mathtt{pii} > ::
    iterator addInterval(set<pii> is, int L, int R) {
  if (L == R) return is.end();
```

# $R = \max(R, \text{ it ->Y});$

```
auto it = is.lower_bound(\{L, R\}), before = it; while (it != is.end() && it->X <= R) {
  before = it = is.erase(it);
if (it != is.begin() && (--it)->Y>= L) {
  L = \min(L, it ->X);
  R = \max(R, it ->Y);
  is.erase(it);
```

```
void removeInterval(set<pii> is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it ->Y;
  if (it->X == L) is.erase(it);
  else (int\&)it ->Y = L;
  if (R != r2) is .emplace(R, r2);
```

return is.insert(before, pii(L, R));

#### 3.4 Leftist Tree [bbd228]

```
11 v, data, sz, sum;
node *1, *r;
  node(ll k)
       v(0), data(k), sz(1), l(0), r(0), sum(k) {}
ll sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
11 node *merge(node *a, node *b) {
  if (!a | | '!b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
  delete tmp;
```

#### 3.5 Heavy light Decomposition\* [babe8a]

```
struct Heavy_light_Decomposition {
  int n, ulink [N], deep [N], mxson [N], w[N], pa [N];
  int t, pl[N], data[N], val[N]; // val: vertex data
```

```
National Taiwan University 8BQube
   \begin{array}{l} {\rm vector}{<} {\rm int}{>} \; G[N] \, ; \\ {\rm void \;\; init} \; ({\rm int} \;\; \underline{\ \ \ \ \ } ) \;\; \{ \end{array}
     n = n;
for (int i = 1; i <= n; ++i)
        G[i].clear(), mxson[i] = 0;
   void add_edge(int a, int b) {
     G[a].pb(b), G[b].pb(a);
   void dfs(int u, int f, int d) {
     w[u] = 1, pa[u] = f, deep[u] = d++; for (int &i : G[u])
         if (i != f) {
           dfs(i, u, d), w[u] += w[i];
            if (w[mxson[u]] < w[i]) mxson[u] = i;
    \begin{array}{lll} void & cut(int \ u, \ int \ link) \ \{ \\ & data[pl[u] = +\!\!+\!\!t] = val[u] \ , \ ulink[u] = link \ ; \end{array} 
      if (!mxson[u]) return;
                                                                                      }
      cut(mxson[u], link);
for (int i : G[u])
                                                                                   };
         if (i != pa[u] && i != mxson[u])
           cut(i, i);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
      \begin{array}{ll} \mbox{int } ta = \mbox{ulink} [a] \,, \ tb = \mbox{ulink} [b] \,, \ res = 0; \\ \mbox{while} \ (ta := tb) \, \{ \end{array}
         if (deep
          [ta] > deep[tb]) swap(ta, tb), swap(a, b);

/ query(pl[tb], pl[b])
         tb = ulink[b = pa[tb]];
      if (pl[a] > pl[b]) swap(a, b);
      // query(pl[a], pl[b])
        Centroid Decomposition* [4eccaf]
void init(int _n) {
     n = \underline{n}, layer[0] =
                                 -1;
      fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
      for (int i = 1; i \le n; ++i) G[i]. clear();
```

```
void add_edge(int a, int b, int w) {
   G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
void get_cent(
  int u, int f, int &mx, int &c, int num) {
   int mxsz = 0;
   sz[u] = 1;
   for (pll e : G[u])
      if (!done[e.X] && e.X != f) {
  get_cent(e.X, u, mx, c, num);
         sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
   if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
void dfs(int u, int f, ll d, int org) {
  // if required, add self info or climbing info
  dis[layer[org]][u] = d;
  for (pll e : G[u])
    if (!done[e.X] && e.X != f)
        dfs(e.X, u, d + e.Y, org);
         dfs(e.X, u, d + e.Y, org);
int cut(int u, int f, int num) {
   int mx = 1e9, c = 0, lc;
get_cent(u, f, mx, c, num);
   done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1; for (pll e : G[c])
      if (!done[e.X]) {
         if (sz[e.X] > sz[c])
  lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
         upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
   return done[c] = 0, c;
```

```
void build() { cut(1, 0, n); }
void modify(int u) {
   for (int a = u, ly = layer[a]; a;
        a = pa[a], --ly) {
        info[a].X += dis[ly][u], ++info[a].Y;
        if (pa[a])
            upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
}

ll query(int u) {
    ll rt = 0;
   for (int a = u, ly = layer[a]; a;
        a = pa[a], --ly) {
        rt += info[a].X + info[a].Y * dis[ly][u];
        if (pa[a])
        rt -=
            upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    }
   return rt;
}
```

#### 3.7 LiChaoST\* [4a61ec]

```
struct L {
  ll m, k, id;
  \begin{array}{l} L() \ : \ id(-1) \ \{\} \\ L(11 \ a, \ 11 \ b, \ 11 \ c) \ : \ m(a) \, , \ k(b) \, , \ id(c) \ \{\} \\ 11 \ at(11 \ x) \ \{ \ return \ m \ ^* \ x + k; \ \} \end{array}
class LiChao { // maintain max
private:
  int n; vector<L> nodes;
  if (nodes[rt].id = -1)
        return nodes[rt] = ln, void();
     bool atLeft = nodes[rt].at(1) < ln.at(1);
     if (nodes[rt].at(m) < ln.at(m))
     atLeft = 1, swap(nodes[rt], ln);
if (r - l == 1) return;
      if \ (atLeft) \ insert(l\,,\,m,\ rt << 1\,,\,ln); \\
     else insert(m, r, rt \ll 1 | 1, ln);
  Il query(int 1, int r, int rt, ll x) {
  int m = (l + r) >> 1; ll ret = -INF;
  if (nodes[rt].id != -1) ret = nodes[rt].at(x);
  if (r - l == 1) return ret;
     public:
  LiChao(int n_) : n(n_), nodes(n * 4) {}
  void insert(\overline{L} ln) { insert(0, n, 1, ln);
  11 query(11 x) { return query(0, n, 1, x); }
```

#### 3.8 Link cut tree\* [5f036a]

```
struct Splay { // xor-sum
  static Splay nil;
  Splay *ch[2], *f;
   int val, sum, rev, size;
   Splay (int
   val = 0) : val(val), sum(val), rev(0), size(1) { f = ch[0] = ch[1] = &nil; }
   bool isr()
   { return f->ch[0] != this && f->ch[1] != this; }
   int dir()
   { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
     ch[d] = c;
      if (c != \&nil) c -> f = this;
     pull();
   void give_tag(int r) {
     if (r) swap(ch[0], ch[1]), rev = 1;
   void push() {
     if (ch[0] != &nil) ch[0]->give_tag(rev);
if (ch[1] != &nil) ch[1]->give_tag(rev);
     rev = 0;
   void pull() {
      // take care of the nil!
      size = ch[0] - size + ch[1] - size + 1;
```

function<bool(const point &, const point &)> f =
 [dep](const point &a, const point &b) {
 if (dep & 1) return a.x < b.x;</pre>

```
else return a.y < b.y;
     if (\operatorname{ch}[1] != \& \operatorname{nil}) \operatorname{ch}[1] -> f = \operatorname{this};
                                                                              int m = (l + r) >> 1;

nth\_element(p + l, p + m, p + r, f);
                                                                              xl[m] = xr[m] = p[m].x;
} Splay::nil;
Splay : nil;
Splay *nil = &Splay :: nil;
void rotate(Splay *x) {
                                                                              yl[m] = yr[m] = p[m] \cdot y;

lc[m] = build(1, m, dep + 1);
void rotate (Splay
  Splay *p = x - > f;
                                                                              if (~lc[m]) {
  int d = x - sdir();
                                                                                 xl[m] = min(xl[m], xl[lc[m]]);
                                                                                 \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{lc}[m]]);
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
p->setCh(x->ch[!d], d);
                                                                                yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
  x->setCh(p, !d);
  p->pull(), x->pull();
                                                                              rc[m] = build(m + 1, r, dep + 1);
                                                                              if (~rc[m]) {
void splay (Splay *x) {
                                                                                xl[m] = min(xl[m], xl[rc[m]]);
  vector < Splay * splay Vec;
for (Splay * q = x;; q = q->f) {
                                                                                 xr[m] = max(xr[m], xr[rc[m]]);
                                                                                 yl[m] = min(yl[m], yl[rc[m]]);
     splayVec.pb(q);
                                                                                yr[m] = max(yr[m], yr[rc[m]]);
     if (q->isr()) break;
                                                                              return m;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
while (!x->isr()) {
                                                                           bool bound (const point &q, int o, long long d) {
                                                                               \begin{array}{l} \mbox{double ds} = \mbox{sqrt}(d+1.0); \\ \mbox{if } (q.x < xl[o] - ds \mid\mid q.x > xr[o] + ds \mid\mid \\ \mbox{q.y} < yl[o] - ds \mid\mid q.y > yr[o] + ds) \end{array} 
     if (x->f->isr()) rotate(x);
     else if (x->dir() = x->f->dir())
       rotate(x->f), rotate(x);
                                                                                 return false;
     else rotate(x), rotate(x);
                                                                              return true;
                                                                           long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
Splay* access(Splay *x) {
  Splay *q = nil;
for (; x != nil; x = x->f)
    \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
                                                                              const point &q, long long &d, int o, int dep = 0) {
  return q;
                                                                              if (!bound(q, o, d)) return;
long long cd = dist(p[o], q);
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
                                                                              if (cd != 0) d = min(d, cd);
  root_path(x), x->give_tag(1);
                                                                              if ((\text{dep \& 1}) \&\& q.x < p[o].x | |
                                                                                !(dep & 1) && q.y < p[o].y) {
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  x->push(), x->pull();
void split (Splay *x, Splay *y) {
  chroot(x), root_path(y);
                                                                              } else {
                                                                                 if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void link (Splay *x, Splay *y) {
  root_path(x), chroot(y);
x->setCh(y, 1);
                                                                              }
                                                                           void init(const vector<point> &v) {
                                                                              for (int i = 0; i < v.size(); ++i) p[i] = v[i];
void cut(Splay *x, Splay *y) {
                                                                              root = build(0, v.size());
  split(x, y);
  if (y->size != 5) return;
  y->push();
                                                                           long long nearest (const point &q) {
  y->ch[0] = y->ch[0]->f = nil;
                                                                              long long res = 1e18;
                                                                              dfs(q, res, root);
Splay* get_root(Splay *x)
                                                                              return res;
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
    x \rightarrow push();
                                                                              // namespace kdt
  splay(x);
                                                                                 Flow/Matching
  return x;
                                                                           4.1 Bipartite Matching* [f07280]
bool conn(Splay *x, Splay *y) {
  return get_root(x) == get_root(y);
                                                                           Splay* lca(Splay *x, Splay *y) {
   access(x), root_path(y);
                                                                              vector < int > G[N + 1];
                                                                              bool dfs(int u) {
  if (y->f = nil) return y;
                                                                                 for (int &i = cur[u]; i < SZ(G[u]); ++i) {
  return y->f;
                                                                                   int e = G[u][i];
                                                                                   if (mq[e]
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
                                                                                          (\operatorname{dis}[\operatorname{mq}[e]] = \operatorname{dis}[u] + 1 \& \operatorname{dfs}(\operatorname{mq}[e]))
                                                                                      return mp[mq[e] = u] = e, 1;
int query (Splay *x, Splay *y) {
                                                                                 return dis[u] = -1, 0;
  split(x, y);
  return y->sum;
                                                                              bool bfs() {
                                                                                 queue<int> q;
                                                                                 fill_n(dis, l + 1, -1);
3.9 KDTree [74016d]
                                                                                 for (int i = 0; i < l; ++i)
                                                                                   if (!~mp[i])
namespace kdt {
int root, lc [maxn], rc [maxn], xl [maxn], xr [maxn],
                                                                                     q.push(i), dis[i] = 0;
                                                                                 while (!q.empty()) {
  yl[maxn], yr[maxn];
point p[maxn];
                                                                                   int u = q.front();
int build(int l, int r, int dep = 0) {
                                                                                   q.pop();
   if (l = r) return -1;
```

if (!~dis[mq[e]])

q.push(mq[e]), dis[mq[e]] = dis[u] + 1;

```
return dis[1] != -1;
}
int matching() {
  int res = 0;
  fill_n(mp, 1, -1), fill_n(mq, r, 1);
  while (bfs()) {
    fill_n(cur, 1, 0);
    for (int i = 0; i < 1; ++i)
        res += (!~mp[i] && dfs(i));
    }
  return res; // (i, mp[i] != -1)
}
void add_edge(int s, int t) { G[s].pb(t); }
void init(int_l, int_r) {
    l = _l, r = _r;
    for (int i = 0; i <= 1; ++i)
        G[i].clear();
}
}</pre>
```

#### 4.2 Kuhn Munkres\* [edf909]

```
{\tt struct} KM { // 0-base, maximum matching
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)
      fill_n (w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
  if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)
         for (int x = 0, y = qu[ql++]; x < n; ++x)
           if (!vl[x] && slk
                [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d; else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x])' d = slk[x];
       for (int x = 0; x < n; ++x) {
         \begin{array}{l} \text{if } (vl[x]) \ hl[x] += d; \\ \text{else } slk[x] -= d; \end{array}
         if (\operatorname{vr}[x]) \operatorname{hr}[x] = d;
       for (int x = 0; x < n; ++x)
         if (!vl[x] && !slk[x] && !Check(x)) return;
  ll solve() {
    fill_n(fl
         , n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
    for (int i = 0; i < n; ++i)
    hl[i] = *max\_element(w[i], w[i] + n);
for (int i = 0; i < n; ++i) bfs(i);
    ll res = 0;
    return res;
};
```

#### 4.3 MincostMaxflow\* [47d9d2]

```
struct MinCostMaxFlow { // 0-base
    struct Edge {
        1l from, to, cap, flow, cost, rev;
    } *past[N];
    vector<Edge> G[N];
    int inq[N], n, s, t;
    ll dis[N], up[N], pot[N];
    bool BellmanFord() {
        fill_n(dis, n, INF), fill_n(inq, n, 0);
    }
}
```

```
queue<int> q;
auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if \ (cap > 0 \ \&\& \ dis [u] > d) \ \{ \\
       dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
  relax(s, 0, INF, 0);
   while (!q.empty())
    int u = q.front();
     q.pop(), inq[u] = 0;
for (auto &e : G[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
             (e.to, d2, min(up[u], e.cap - e.flow), &e);
    }
  }
  return dis[t] != INF;
void solve (int
  , int _t, ll &flow, ll &cost, bool neg = true) { s = \_s, t = \_t, flow = 0, cost = 0;
   if (neg) BellmanFord(), copy_n(dis, n, pot);
   for (; BellmanFord(); copy_n(dis, n, pot)) {
     for (int
     i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
     for (int i = t; past[i]; i = past[i]->from) {
  auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
  }
}
void init(int _n) {
    n = _n, fill_n(pot, n, 0);
  for (int i = 0; i < n; ++i) G[i].clear();
void add_edge(ll a, ll b, ll cap, ll cost)
  \begin{array}{l} G[a].p\overline{b}(Edge\{a,\ b,\ cap\ ,\ 0,\ cost\ ,\ SZ(G[b])\ ))\,;\\ G[b].pb(Edge\{b,\ a,\ 0,\ 0,\ -cost\ ,\ SZ(G[a])\ -\ 1\}); \end{array}
```

# 4.4 Maximum Simple Graph Matching\* [233755]

```
struct Matching { // 0-base
   queue < int > q; int n;
    vector<int> fa, s, vis, pre, match;
    vector<vector<int>>> G;
   int Find(int u)
    \{ return u = fa[u] ? u : fa[u] = Find(fa[u]); \}
    int LCA(int x, int y) {
      static int tk = 0; tk++; x = Find(x); y = Find(y); for (;; swap(x, y)) if (x != n) {
    if (vis[x] == tk) return x;
         vis[x] = tk;
         x = Find(pre[match[x]]);
   void Blossom(int x, int y, int l) {
  for (; Find(x) != l; x = pre[y]) {
         pre[x] = y, y = match[x];
         if (s[y] = 1) q.push(y), s[y] = 0;
for (int z: \{x, y\}) if (fa[z] = z) fa[z] = 1;
   bool Bfs(int r) {
      \begin{array}{ll} iota\left(ALL(fa),\ 0\right); & fill\left(ALL(s),\ -1\right); \\ q = queue < int > (); & q.push(r); & s[r] = 0; \end{array}
       for (; !q.empty(); q.pop()) {
          for (int x = q.front(); int u : G[x])
             if(s[u] = -1) {
                if (pre[u] = x, s[u] = 1, match[u] == n) {
                   for (int a = u, b = x, last;
                         b != n; a = last, b = pre[a]
                      last =
                            match[b], match[b] = a, match[a] = b;
                   return true:
              \begin{array}{l} \text{q.push}(\mathrm{match}[u])\,; \;\; s[\mathrm{match}[u]] = 0\,; \\ \} \;\; \text{else} \;\; \text{if} \;\; (!\,s[u] \;\; \&\& \;\; \mathrm{Find}(u) \;\; != \;\; \mathrm{Find}(x)) \;\; \{ \end{array} 
                int l = LCA(u, x);
Blossom(x, u, l); Blossom(u, x, l);
```

return false;

```
National Taiwan University 8BQube
   Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis
   \begin{array}{c} (n+1)\,,\; pre\,(n+1,\;n)\,,\; match\,(n+1,\;n)\,,\; G(n)\ \{\}\\ void\ add\_edge\,(int\ u,\; int\ v) \end{array}
    \left\{ \begin{array}{l} G[u].pb(v), G[v].pb(u); \\ int \ solve() \end{array} \right\} 
      int ans = 0;
      for (int x = 0; x < n; ++x)
         if (match[x] == n) ans += Bfs(x);
      return ans;
  } // match[x] = n means not matched
       Maximum Weight Matching* [c80005]
#define REP(i, l, r) for (int i=(l); i<=(r); ++i) struct WeightGraph { // 1-based
   struct edge { int u, v, w; }; int n, nx;
   vector<int> lab; vector<vector<edge>>> g;
   vector<int> slk, match, st, pa, S, vis;
  vector<int> six, match, st, pa, b, vis, vector<vector<int> flo, flo_from; queue<int> q; WeightGraph(int n_): n(n_), nx(n * 2), lab(nx + 1), g(nx + 1, vector<edge>(nx + 1)), slk(nx + 1),
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
      match = st = pa = S = vis = slk;
     \label{eq:rep} \text{REP}(u, \ 1, \ n) \ \hat{\text{REP}}(v, \ 1, \ n) \ g[u][v] = \{u, \ v, \ 0\};
   int E(edge e)
   { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; } void update_slk(int u, int x, int &s)
   { if (!s | E(g[u][x]) < E(g[s][x]))'s = u; } void set_slk(int x) {
     slk[x] = 0;

REP(u, 1, n)
         if (g[u][x].w > 0 \& st[u] != x \& S[st[u]] == 0)
            update_slk(u, x, slk[x]);
   void q_push(int x)
      if (x \le n) q.push(x);
      else for (int y : flo[x]) q_push(y);
   void set_st(int x, int b) {
      st[x] = b;
      if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
         reverse(1 + ALL(f)), it = f.end() - pr;
      auto res = vector(f.begin(), it);
                                                                                                       }
      return f.erase(f.begin(), it), res;
   void set_match(int u, int v) {
  match[u] = g[u][v].v;
      if (u \le n) return;
      \begin{array}{l} \text{int } xr = flo\_from[u][g[u][v].u]; \\ \text{auto } \&f = flo[u], \ z = split\_flo(f, \ xr); \\ \text{REP}(i, \ 0, SZ(z) - 1) \ \text{set\_match}(z[i], \ z[i \ \cap \ 1]); \\ \end{array} 
      set_match(xr, v); f.insert(f.end(), ALL(z));
   void augment(int u, int v) {
      for (;;) {
         int xnv = st[match[u]]; set\_match(u, v);
         if (!xnv) return;
         set_match(v = xnv, u = st[pa[xnv]]);
   int lca(int u, int v) {
      static int t = 0; ++\hat{t};
      for (++t; u | | v; swap(u, v)) if (u) {
  if (vis[u] == t) return u;
  vis[u] = t, u = st[match[u]];
         if (u) u = st[pa[u]];
      return 0;
   void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0, match[b] = match[o];
      vector < int > f = \{o\};
```

for (int t : {u, v}) {
 reverse(1 + ALL(f));

 $flo[b] = f; set\_st(b, b);$ 

for (int x = t, y; x != o; x = st[pa[y]])

 $REP(x\,,\ 1\,,\ nx)\ g\,[\,b\,]\,[\,x\,]\,.w\,=\,g\,[\,x\,]\,[\,b\,]\,.w\,=\,0\,;$ 

f.pb(x), f.pb(y = st[match[x]]),  $q_push(y)$ ;

```
\begin{array}{l} \label{eq:condition} \text{fill} \left( \text{ALL} \left( \text{flo\_from} \left[ b \right] \right) \,, \,\, 0 \right); \\ \text{for} \,\, \left( \text{int} \,\, xs \,: \,\, \text{flo} \left[ b \right] \right) \,\, \left\{ \end{array}
          g[b][x] = g[xs][x], g[x][b] = g[x][xs];
                              1, n)
                 if (flo\_from[xs][x]) flo\_from[b][x] = xs;
     set_slk(b);
void expand_blossom(int b) {
     for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
     for (int x: split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
           slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
      for (int x : flo[b])
            if (x = xr) S[x] = 1, pa[x] = pa[b];
            else S[x] = -1, set\_slk(x);
      st[b] = 0;
bool on_found_edge(const edge &e) {
      \begin{array}{lll} & \text{if (int } u = st \, [e \cdot u] \;, \; v = st \, [e \cdot v] \;; \; S[v] = -1) \; \{ \\ & \text{int } nu = st \, [match[v]] \;; \; pa[v] = e \cdot u \;; \; S[v] = 1 \;; \\ & slk \, [v] = slk \, [nu] = S[nu] = 0 \;; \; q\_push(nu) \;; \\ \end{array} 
      else if (S[v] = 0)
            if (int o = lca(u, v)) add_blossom(u, o, v);
           else return augment(u, v), augment(v, u), true;
     return false;
bool matching() { fill(ALL(S), -1), fill(ALL(slk), 0);
      q = queue < int > ();
     \begin{array}{l} \text{REP}(x, 1, nx) & \text{if } (st[x] = x \& x : match[x]) \\ pa[x] = S[x] = 0, \text{ q_push}(x); \end{array}
       if (q.empty()) return false;
      for (;;)
            while (SZ(q)) {
                 int u = q.front(); q.pop();
                  if (S[st[u]] = 1) continue;
                REP(v, 1, n)
                       if (g[u][v].w > 0 & st[u] != st[v]) {
                             if (E(g[u][v]) != 0)
                                  update\_slk(u, st[v], slk[st[v]]);
                                            (on_found_edge(g[u][v])) return true;
            int d = INF;
          REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
           REP(x,\ 1,\ nx)
                 if (int
                      s = slk[x]; st[x] == x && s && s[x] <= 0)

d = min(d, E(g[s][x]) / (S[x] + 2));
           REP(u, 1, n)
                 if (S[st[u]] = 1) lab[u] += d;
                 else if (S[st[u]] = 0) {
   if (lab[u] \le d) return false;
                       lab[u] -= d;
          REP(x, 1, nx)
                 if (int s = slk[x]; st[x] == x &&
                       \begin{array}{c} s \&\& st[s] \mathrel{!=} x \&\& E(g[s][x]) == 0) \\ if (on\_found\_edge(g[s][x])) \ return \ true; \end{array}
           REP(b\,,\ n\,+\,1\,,\ nx\,)
                  if (st[b] = b \&\& S[b] = 1 \&\& lab[b] = 0)
                      expand_blossom(b);
     return false;
pair < ll, int > solve() {
      fill (ALL(match), 0);
      \begin{array}{lll} & \text{REP}(\mathbf{u}, \ \mathbf{0}, \ \mathbf{n}) \ \ \text{st} \ [\mathbf{u}] = \mathbf{u}, \ \ \text{flo} \ [\mathbf{u}]. \ \text{clear} \ () \ ; \\ & \text{int} \ \ \mathbf{w} \_ \text{max} = \ \mathbf{0} \ ; \\ & \end{array} 
      \begin{array}{lll} \text{REP}(u, 1, n) & \text{REP}(v, 1, n) & \{ & \\ \text{flo\_from} \left[u\right] \left[v\right] & = \left(u = v ? u : 0\right); \end{array} 
           w_m = max(w_m x, g[u][v].w);
      fill (ALL(lab), w_max);
     \begin{tabular}{ll} \beg
```

return 0;

bool bfs() {

 $fill_n(dis, n + 3, -1);$ 

```
\begin{array}{l} \text{queue} < \text{int} > \text{q}; \\ \text{q.push}(\text{s}), & \text{dis}[\text{s}] = 0; \end{array}
      while (matching()) ++n_matches;
     REP(u, 1, n) if (match[u] \&\& match[u] < u)
        tot_weight += g[u][match[u]].w;
                                                                                      while (!q.empty()) {
     return make_pair(tot_weight, n_matches);
                                                                                         int u = q. front();
                                                                                         q.pop();
  void add_edge(int u, int v, int w)
{ g[u][v].w = g[v][u].w = w; }
                                                                                         for (edge &e : G[u])
if (!~dis[e.to] && e.flow != e.cap)
                                                                                              q.push(e.to), dis[e.to] = dis[u] + 1;
4.6 SW-mincut [90bfe6]
                                                                                      return dis[t] != -1;
int maxflow(int _s, int _t) {
                                                                                      s = \_s, t = \_t;
                                                                                      int \overline{flow} = 0, df;
  \begin{array}{ll} \text{int vst} \left[ \text{MXN} \right], \text{ edge} \left[ \text{MXN} \right] \left[ \text{MXN} \right], \text{ wei} \left[ \text{MXN} \right]; \\ \text{void init} \left( \text{int n} \right) \end{array} \right\}
                                                                                      while (bfs()) {
                                                                                         fill_n(cur, n + 3, 0);
while ((df = dfs(s, INF))) flow += df;
     REP fill_n(edge[i], n, 0);
   void addEdge(int u, int v, int w){
                                                                                      return flow;
     edge[u][v] += w; edge[v][u] += w;
                                                                                   bool solve() {
   int search(int &s, int &t, int n){
                                                                                      int sum = 0;
     fill_n(vst, n, 0), fill_n(wei, n, 0);
                                                                                      for (int i = 0; i < n; ++i)
     s = t = -1;
                                                                                         if (cnt[i] > 0)
     int mx, cur;
                                                                                      add_edge(n + 1, i, cnt[i]), sum += cnt[i];
else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);
if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int j = 0; j < n; +++j) {
       mx = -1, cur = 0;

mx = -1, cur = 0;

mx = -1, mx = wei[i];
                                                                                      for (int i = 0; i < n; ++i)
        vst[cur] = 1, wei[cur] = -1;
                                                                                             (\,\mathrm{cnt}\,[\,i\,]\,>\,0)
        s = t; t = cur;
                                                                                         G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)
        G[i].pop\_back(), G[n + 2].pop\_back();
     return mx;
                                                                                      return sum != -1;
   int solve(int n) {
                                                                                   int solve(int _s, int
                                                                                                                  t) {
     int res = INF;
                                                                                      add_edge(_t, _s, INF);
if (!solve()) return -1; // invalid flow
     for (int x, y; n > 1; n--){
         \begin{array}{l} res = min(res, search(x, y, n)); \\ REP \ edge[i][x] = (edge[x][i] += edge[y][i]); \\ \end{array} 
                                                                                      int x = G[\underline{t}] . back() . flow;
                                                                                      return G[_t].pop_back(), G[_s].pop_back(), x;
        REP {
          edge[y][i] = edge[n - 1][i];
edge[i][y] = edge[i][n - 1];
                                                                                };
          // edge[y][y] = 0; 
                                                                                4.8 Gomory Hu tree* [5f2460]
                                                                                MaxFlow Dinic;
     return res;
                                                                                int g[MAXN];
  }
} sw;
                                                                                void GomoryHu(int n) { // 0-base
                                                                                   fill_n(g, n, 0);
for (int i = 1; i < n; ++i) {
4.7
       BoundedFlow*(Dinic*) [4ae8ab]
                                                                                      Dinic.reset();
{\tt struct} \;\; {\tt BoundedFlow} \; \left\{ \;\; // \;\; {\tt 0-base} \right.
                                                                                      add\_edge(\,i\;,\;g\,[\,i\,]\;,\;Dinic.maxflow(\,i\;,\;g\,[\,i\,]\,)\,)\,;
   struct edge {
                                                                                      for (int j = i + 1; j \le n; ++j)

if (g[j] = g[i] \&\& \sim Dinic.dis[j])
     int to, cap, flow, rev;
                                                                                           g[j] = i;
   vector<edge> G[N];
   int n, s, t, dis [N], cur [N], cnt [N];
                                                                                }
   void init(int _n) {
     n = _n;
for (int i = 0; i < n + 2; ++i)
                                                                                       Minimum Cost Circulation* [cb40c6]
        G[i].clear(), cnt[i] = 0;
                                                                                struct MinCostCirculation { // 0-base
                                                                                   {\tt struct} \  \, {\tt Edge} \  \, \{
   void add_edge(int u, int v, int lcap, int rcap) {
                                                                                      ll from, to, cap, fcap, flow, cost, rev;
     } *past[N];
                                                                                    vector<Edge> G[N];
                                                                                   ll dis [N], inq [N], n;
void BellmanFord(int s) {
  void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, SZ(G[v])});
  G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
                                                                                      fill_n(dis, n, INF), fill_n(inq, n, 0);
                                                                                      queue<int> q;
                                                                                      auto relax = [&](int u, ll d, Edge *e) {
                                                                                         if (dis[u] > d) {
   int dfs(int u, int cap) {
  if (u = t || !cap) return cap;
                                                                                           \begin{array}{l} \text{dis}\,[u] = d, \;\; \text{past}\,[u] = e; \\ \text{if} \;\; (!\, \text{inq}\,[u]) \;\; \text{inq}\,[u] = 1, \;\; \text{q.push}(u); \end{array}
     for (int &i = cur[u]; i < SZ(G[u]); ++i) {
                                                                                        }
        edge \&e = G[u][i];
        if (dis[e.to] == dis[u] + 1 && e.cap!= e.flow) {
                                                                                      relax(s, 0, 0);
           int df = dfs(e.to, min(e.cap - e.flow, cap));
                                                                                      while (!q.empty()) {
           if (df) {
                                                                                         int u = q.front();
              e.flow += df, G[e.to][e.rev].flow -= df;
                                                                                         q.pop(), inq[u] = 0;
              return df;
                                                                                         for (auto &e : G[u])
          }
                                                                                           if(e.cap > e.flow)
                                                                                              relax\left(\begin{smallmatrix}e.to\end{smallmatrix}, \; dis\left[\begin{matrix}u\end{smallmatrix}\right] \; + \; e.cost\;, \; \&e\right);
        }
     dis[u] = -1;
```

void try\_edge(Edge &cur) {

if (dis[cur.from] + cur.cost < 0) {

BellmanFord(cur.to);

if (cur.cap > cur.flow) return ++cur.cap, void();

#### 4.10 Flow Models

- $\bullet \quad {\rm Maximum/Minimum\,flow\,with\,lower\,bound/Circulation\,problem}$ 
  - 1. Construct supersource S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - − To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$  , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite  $\operatorname{graph}(X,Y)$ 
  - 1. Redirect every edge:  $y \to x$  if  $(x,y) \in M, x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- $\bullet \quad {\rm Minimum} \, {\rm cost} \, {\rm cyclic} \, {\rm flow}$ 
  - 1. Consruct super source S and  $\sinh T$
  - 2. For each edge (x, y, c), connect  $x \to y$  with (cost, cap) = (c, 1) if c > 0, otherwise connect  $y \to x$  with (cost, cap) = (-c, 1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - $4. \ \ \text{For each vertex} \ v \ \text{with} \ d(v) > 0, \\ \text{connect} \ S \rightarrow v \ \text{with} \ (cost, cap) = (0, d(v))$
  - 5. For each vertex v with d(v) < 0, connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source  $s \to v, v \in G$  with capacity K
- 4. For each edge (u, v, w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
- 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow

- 1. Capacity  $c_{uv}$ , Flow  $f_{uv}$ , Cost  $w_{uv}$ , Required Flow difference for vertex  $b_u$ .
- 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$
  $p_{u} \ge 0$ 

## 5 String

#### 5.1 KMP [9e1cd1]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
}
for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
}
return ans;
}</pre>
```

#### **5.2 Z-value\*** [e2dc6f]

### 5.3 Manacher\* [bfe74e]

### **5.4 SAIS\*** [e9a275]

```
auto sais (const auto &s) {
  const int n = SZ(s), z = ranges::max(s) + 1;
  if (n = 1) return vector\{0\};
vector<int> c(z); for (int x : s) ++c[x];
  partial_sum(ALL(c), begin(c));
  vector < int > sa(n); auto I = views :: iota(0, n);
  vector<br/>bool> t(n, true);
  for (int i = n - 2; i \ge 0; --i)
    t[i] = (
         s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
  auto is_lms = views::filter([&t](int x) {
    return x && t[x] && !t[x - 1];
  });
  auto induce = [&] {
    for (auto x = c; int y : sa)

if (y--) if (!t[y]) sa[x[s[y]-1]++] = y;

for (auto x = c; int y : sa | views::reverse)
       if (y--) if (t[y]) sa[--x[s[y]]] = y;
  vector < int > lms, q(n); lms.reserve(n);
  for (auto x = c; int i : I | is_lms)
    q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
  induce(); vector<int> ns(SZ(lms));
```

for (int j = -1, nz = 0; int  $i : sa \mid is_lms$ ) {

```
if (j >= 0) {
        int len = min({n - i, n - j, lms[q[i] + 1] - i});
        ns[q[i]] = nz += lexicographical_compare(
             begin(s) + j, begin(s) + j + len
             begin(s) + i, begin(s) + i + len);
     j \; = \; i \; ;
   fill(ALL(sa), 0); auto nsa = sais(ns);
  for (auto x = c; int y : nsa | views::reverse)
     y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
// sa[i]: sa[i]-th suffix
       is the i-th lexicographically smallest suffix.
   hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
  int n; vector<int> sa, hi, ra;
  Suffix
     (const auto &_s, int _n) : n(_n), hi(n), ra(n) { vector<int> s(n + 1); // s[n] = 0; copy_n(_s, n, begin(s)); // _s shouldn't contain 0
     sa = sais(s); sa.erase(sa.begin());
     for (int i = 0; i < n; ++i) ra[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {
  if (!ra[i]) { h = 0; continue; }
  for (int j = sa[ra[i] - 1]; max
        (i, j) + h < n & & s[i + h] == s[j + h];) ++h;

hi[ra[i]] = h ? h-- : 0;
  }
};
```

#### 5.5 Aho-Corasick Automatan\* [91c6c0]

```
struct AC_Automatan {
   int nx[len][sigma], fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
   int newnode() {
      fill_n(nx[top], sigma, -1);
      return top++;
   void init() { top = 1, newnode(); }
   int input(string &s) {
      int X = 1;
      \begin{array}{lll} & \text{for } (c\text{har } c:s) \; \{ & \text{if } (!\!\sim\!\!\operatorname{nx}[X][c-'A']) \; \operatorname{nx}[X][c-'A'] = \operatorname{newnode}() \, ; \\ & X = \operatorname{nx}[X][c-'A'] \, ; \end{array}
      return X; // return the end node of string
   void make_fl() {
      queue\langle int \rangle q;
q.push(1), fl[1] = 0;
       for (int t = 0; !q.empty(); ) {
          int R = q.front();
          q.pop(), ord[t++] = R;
          for (int i = 0; i < sigma; ++i)
if (~nx[R][i]) {
                int X = rnx[R][i] = nx[R][i], Z = fl[R];
                 for (; Z \&\& !\sim nx[Z][i];) Z = fl[Z];

fl[X] = Z ? nx[Z][i] : 1, q.push(X);
             else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
      }
   void solve() {
      \begin{array}{lll} & \text{for (int i = top - 2; i > 0; --i)} \\ & \text{cnt[fl[ord[i]]] += cnt[ord[i]];} \end{array}
} ac;
```

#### 5.6 Smallest Rotation [e74dc0]

```
\begin{array}{l} string \ mcp(string \ s) \ \{\\ int \ n = SZ(s) \,, \ i = 0 \,, \ j = 1;\\ s += s \,;\\ while \ (i < n \&\& \ j < n) \ \{\\ int \ k = 0 \,;\\ while \ (k < n \&\& \ s[i + k] \Longrightarrow s[j + k]) \ +\!\!\!+k;\\ if \ (s[i + k] <= s[j + k]) \ j +\!\!\!= k + 1;\\ else \ i +\!\!\!= k + 1;\\ if \ (i \Longrightarrow j) +\!\!\!+j;\\ \}\\ int \ ans = i < n \ ? \ i : j;\\ return \ s.substr(ans, n);\\ \} \end{array}
```

### 5.7 De Bruijn sequence\* [f601c2]

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N
  void dfs(int *out, int t, int p, int &ptr) {
    if (ptr >= L) return;
    if (t > N) {
        if (N % p) return;
        for (int i = 1; i <= p && ptr < L; ++i)
            out[ptr++] = buf[i];
    } else {
        buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
        for (int j = buf[t - p] + 1; j < C; ++j)
            buf[t] = j, dfs(out, t + 1, t, ptr);
        }
    }
    void solve(int _c, int _n, int _k, int *out) {
        int p = 0;
        C = _c, N = _n, K = _k, L = N + K - 1;
        dfs(out, 1, 1, p);
        if (p < L) fill(out + p, out + L, 0);
    }
} dbs;</pre>
```

```
5.8 Extended SAM* [58fa19]
struct exSAM { int len [N * 2], link [N * 2]; // maxlength, suflink int next [N * 2] [CNUM], tot; // [0, tot), root = 0 int lenSorted [N * 2]; // topo. order
   int cnt[N * 2]; // occurence
   int newnode()
     fill_n (next[tot], CNUM, 0);
     len[tot] = cnt[tot] = link[tot] = 0;
     return tot++;
  void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
  int cur = next[last][c];
     len[cur] = len[last] + 1;
int p = link[last];
     while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
     int q = next[p][c];
     if (len
          [p] + 1 \Longrightarrow len[q] return link[cur] = q, cur;
     int clone = newnode();
     for (int i = 0; i < CNUM; ++i)
       next
             [i] = len[next[q][i]] ? next[q][i] : 0;
     len[clone] = len[p] + 1;
     while (p != -1 \&\& next[p][c] == q)
     next[p][c] = clone, p = link[p];
link[link[cur] = clone] = link[q];
     link[q] = clone;
     return cur;
   void insert(const string &s) {
     int cur = 0:
     for (auto ch : s) {
       int & mxt = next[cur][int(ch - 'a')];
        if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
  void build() {
     queue < int > q;
     q.push(0);
     while (!q.empty()) {
       int cur = q.front();
        q.pop();
        for (int i = 0; i < CNUM; ++i)
          if \(\text[cur][i])
            q.push(insertSAM(cur, i));
     vector<int> lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];
     partial_sum(ALL(lc), lc.begin());
     for (int i
          = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  void solve() {
     for (int i = tot - 2; i >= 0; --i)
```

cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];

```
| };
 5.9 PalTree* [675736]
 struct palindromic_tree {
   struct node {
     int next[26], fail, len;
     int cnt, num; // cnt: appear times, num: number of
    // pal. suf.
     node(int l = 0): fail(0), len(1), cnt(0), num(0) {
        for (int i = 0; i < 26; ++i) next[i] = 0;
   vector<node> St;
   vector<char> s;
   int last, n;
   \begin{array}{lll} palindromic\_tree() : St(2), \ last(1), \ n(0) \ \{ \\ St[0].fail = 1, \ St[1].len = -1, \ s.pb(-1); \end{array}
   inline void clear() {
     St.clear(), s.clear(), last = 1, n = 0;
St.pb(0), St.pb(-1);
     St[0].fail = 1, s.pb(-1);
   inline int get_fail(int x) {
     while (s[\tilde{n} - St[x].len - 1] != s[n])
       x = \hat{S}t[x].fail;
     return x;
   inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
     int cur = get_fail(last);
     if (!St[cur].next[c]) {
        int now = SZ(St);
        St.pb(St[cur].len + 2);
        St [now]. fail =
          St[get_fail(St[cur].fail)].next[c];
        St[cur].next[c] = now;
        St[now].num = St[St[now].fail].num + 1;
     last = St[cur].next[c], ++St[last].cnt;
   inline void count() { // counting cnt
     auto i = St.rbegin();
     for (; i != St.rend(); ++i) {
        St[i-stail].cnt = i-scnt;
   inline int size() { // The number of diff. pal.
     return SZ(St) - 2;
 5.10 Main Lorentz [eaf279]
 vector < pair < int, int \gg rep[kN]; // 0-base [l, r]
```

```
void main_lorentz(const string &s, int sft = 0) {
  const int n = s.size();
  if (n = 1) return;
  const int nu = n / 2, nv = n - nu;
  const string u = s.substr(0, nu), v = s.substr(nu),
         ru(u.rbegin
              (), u.rend()), rv(v.rbegin(), v.rend());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u), z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v); auto get_z = [](const vector<int>&z, int i) {
    return
          (0 \le i \text{ and } i < (int)z.size()) ? z[i] : 0; };
  auto add_rep
        =\,[\,\&\,](\,bool\ left\;,\;int\ c\,,\;int\ l\,,\;int\ k1\,,\;int\ k2\,)\ \{
    const
          int L = \max(1, 1 - k2), R = \min(1 - left, k1);
    if (L > R) return;
    if (left)
          rep[1].emplace\_back(sft + c - R, sft + c - L);
    else rep[l].emplace_back
         (sft + c - R - l + 1, sft + c - L - l + 1);
  for (int cntr = 0; cntr < n; cntr++) {
    int 1, k1, k2;
    if (cntr < nu) {
       l = nu - cntr;
      k1 = get_z(z1, nu - cntr);

k2 = get_z(z2, nv + 1 + cntr);
    } else {
       l = cntr - nu + 1;
```

```
 \begin{array}{c} k1 = \operatorname{get}\_z(z3, \ nu + 1 + nv - 1 - (cntr - nu)); \\ k2 = \operatorname{get}\_z(z4, \ (cntr - nu) + 1); \\ \} \\ if \ (k1 + k2 >= 1) \\ add\_rep(cntr < nu, \ cntr, \ l, \ k1, \ k2); \\ \} \\ // \ p \ (in \ [l, \ r] \implies s[p, \ p + i) = s[p + i, \ p + 2i) \end{array}
```

#### 6 Math

#### 6.1 ax+by=gcd(only exgcd \*) [5fef50]

```
\begin{array}{l} pll \ exgcd(ll \ a, \ ll \ b) \ \{\\ if \ (b = 0) \ return \ pll(1, \ 0);\\ ll \ p = a \ / \ b;\\ pll \ q = exgcd(b, \ a \% \ b);\\ return \ pll(q.Y, \ q.X \ - \ q.Y \ ^* \ p);\\ \}\\ /^* \ ax + by = res \ , \ let \ x \ be \ minimum \ non-negative\\ g, \ p = gcd(a, \ b) \ , \ exgcd(a, \ b) \ ^* \ res \ / \ g\\ if \ p.X < 0: \ t = (abs(p.X) + b \ / \ g \ - 1) \ / \ (b \ / \ g)\\ else: \ t = -(p.X \ / \ (b \ / \ g))\\ p \ += \ (b \ / \ g, \ -a \ / \ g) \ ^* \ t \ ^*/\\ \end{array}
```

#### 6.2 Floor and Ceil [1ffa73]

#### 6.3 Floor Enumeration [67ad61]

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
  int x = n / l;
  r = n / x;
}</pre>
```

#### **6.4** Mod Min [038fef]

```
// min{k | l <= ((ak) mod m) <= r}, no solution -> -1
ll mod_min(ll a, ll m, ll l, ll r) {
  if (a == 0) return l ? -1 : 0;
  if (ll k = (l + a - 1) / a; k * a <= r)
    return k;
  ll b = m / a, c = m % a;
  if (ll y = mod_min(c, a, a - r % a, a - l % a))
    return (l + y * c + a - l) / a + y * b;
  return -1;
}</pre>
```

#### 6.5 Gaussian integer gcd [4fcbff]

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
     (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
     ll c = a.real() * b.real() + a.imag() * b.imag();
     ll d = a.imag() * b.real() - a.real() * b.imag();
     ll r = b.real() * b.real() + b.imag() * b.imag();
     if (c % r == 0 && d % r == 0) return b;
     return gaussian_gcd
          (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
}
```

#### **6.6** Miller Rabin\* [14b81a]

#### 6.7 Simultaneous Equations [21b2e1]

```
struct matrix { //m variables, n equations
   int n, m;
   \label{eq:maxn} \textit{fraction} \ \ M[\textit{MAXN}] \left[ \textit{MAXN} + \ 1 \right], \ \ \textit{sol} \left[ \textit{MAXN} \right];
   int solve() { //-1: inconsistent, >= 0: rank
for (int i = 0; i < n; ++i) {</pre>
         int piv = 0;
         while (piv < m && !M[i][piv].n) ++piv;
         if (piv == m) continue;
         for (int j = 0; j < n; +++j) {
            if (i == j) continue;
             \begin{array}{ll} fraction & tmp = -M[\,j\,] \big[\,piv\,] \ / \ M[\,i\,] \big[\,piv\,]\,; \\ for & (int \ k = 0\,; \ k <= \end{array} 
                   m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
         }
      \inf_{int} rank = 0;

\inf_{int} (int i = 0; i < n; ++i) {
         int piv = 0;
         while (piv < m \&\& !M[i][piv].n) ++piv;
         if (piv = m \&\& M[i][m].n) return -1;
         else if (piv
                < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
      return rank;
};
```

#### 6.8 Pollard Rho\* [fff0fc]

#### 6.9 Simplex Algorithm [40618e]

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
\begin{array}{lll} \textbf{double} & \textbf{a} \, [\text{MAXN}] \, [\text{MAXM}] \; , \; \; \textbf{b} \, [\text{MAXN}] \; , \; \; \textbf{c} \, [\text{MAXM}] \; ; \\ \textbf{double} & \textbf{d} \, [\text{MAXM}] \, [\text{MAXM}] \; , \; \; \textbf{x} \, [\text{MAXM}] \; ; \\ \end{array}
int ix[MAXN + MAXM]; // !!! array all indexed from 0
// \max\{cx\} \text{ subject to } \{Ax \le b, x > = 0\}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector // usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
   fill_n(d[n], m + 1, 0);
fill_n(d[n + 1], m + 1, 0);
   iota(ix, ix + n + m, 0);
   int r = n, s = m - 1;
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j]; d[i][m - 1] = 1;
      d[i][m] = b[i];
      if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);

d[n + 1][m - 1] = -1;
   for (double dd;; ) {
      if (r < n) {
        d[i][j] += d[r][j] * d[i][s];
d[i][s] *= d[r][s];
```

```
}
 if'(s < 0) break;
 for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
   if (r < 0 | |
      (dd = d[r][m]
          \frac{1}{d[r][s]} - d[i][m] / d[i][s] < -eps | |
      (dd < eps \&\& ix[r+m] > ix[i+m]))
 if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not executable
double ans = 0:
fill_n(x, m, 0);
for (int i = m; i <
 x[ix[i]] = d[i-m][m];
return ans;
```

#### 6.9.1 Construction

Primal	Dual
Maximize $c^{\intercal}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \ge c, y \ge 0$
Maximize $c^{\intercal}x$ s.t. $Ax \leq b$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y = c, y \ge 0$
Maximize $c^{\intercal}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \geq c$

 $\overline{\mathbf{x}}$  and  $\overline{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\overline{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\overline{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\overline{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\overline{x}_j = b_j$  holds.

```
1. In case of minimization, let c_i' = -c_i

2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \rightarrow \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j

3. \sum_{1 \le i \le n} A_{ji} x_i = b_j

• \sum_{1 \le i \le n} A_{ji} x_i \le b_j

• \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
```

4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i - x_i'$ 

#### 6.10 chineseRemainder [fe9f25]

#### 6.11 Factorial without prime factor\* [dcffcb]

```
// O(p^k + log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
     if (i % p) prod[i] = prod[i - 1] * i % pk;
     else prod[i] = prod[i - 1];
ll rt = 1;
for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
return rt;
} // (n! without factor p) % p^k</pre>
```

#### 6.12 Discrete Log\* [ba4ac0]

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {
      p[y] = i;
      y = 1LL * y * x % m;
      b = 1LL * b * x % m;
}</pre>
```

```
for (int i = 0; i < m + 10; i += kStep) {
s = 1LL * s * b \% m;
    if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m = 1) return 0;
  for (int i = 0; i < 100; ++i) {
    if (s == y) return i;
s = 1LL * s * x % m;
  if (s = y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
```

#### 6.13 Berlekamp Massey [9380b8]

```
template <typename T>
 vector<T> BerlekampMassey(const vector<T> &output) {
                vector < T > d(SZ(output) + 1), me, he;
               for (int f = 0, i = 1; i <= SZ(output); ++i) {
  for (int j = 0; j < SZ(me); ++j)
   d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
                                  if (me.empty()) {
                                              me. resize (\hat{f} = i);
                                               continue;
                            \begin{array}{l} \mbox{vector} < T > o(i - f - 1); \\ T \ k = -d[i] \ / \ d[f]; \ o.pb(-k); \\ \mbox{for} \ (T \ x : he) \ o.pb(x * k); \\ \mbox{o.resize} (max(SZ(o), SZ(me))); \\ \mbox{for} \ (int \ j = 0; \ j < SZ(me); ++j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); ++j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); ++j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); ++j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j = 0; \ j < SZ(me); +-j) \ o[j] += me[j]; \\ \mbox{if} \ (int \ j =
                                if (i - f + SZ(he)) = SZ(me) he = me, f = i;
                              me = o:
                return me:
```

#### 6.14 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679\ 999983\ 1097774749\ 1076767633\ 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 \ 98789101 \ 987777733 \ 999991921 \ 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 \ \ 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

#### 6.15 Theorem

• Cramer's rule

$$ax+by=e \Rightarrow x = \frac{ed-bf}{ad-bc}$$

$$cx+dy=f \Rightarrow y = \frac{af-ec}{ad-bc}$$

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

• Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i), L_{ij} = -c$  where cisthenumber of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the  $\max \operatorname{maximum} \operatorname{matching} \operatorname{on} G.$ 

- Cayley's Formula
  - Given a degree sequence  $d_1, d_2, ..., d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex 1,2,...,k belong to different components. Then  $T_{n,k}\!=\!kn^{n-k-1}$  .
- Erdős–Gallaitheorem

A sequence of nonnegative integers  $d_1 \ge \cdots \ge d_n$  can be represented as the  ${\tt degree sequence of a finite simple graphon} n vertices if and only if d_1 + \dots + d_n$ 

is even and 
$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i,k)$$
 holds for every  $1 \le k \le n$ .

• Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \ldots, b_n$  is bigraphic if and only if  $\sum_{i=1}^{n} a_i = \sum_{i=1}^{n} b_i$  and  $\sum_{i=1}^{k} a_i \leq \sum_{i=1}^{n} \min(b_i, k)$  holds for

Fulkerson–Chen–Ansteetheorem A sequence  $(a_1, b_1), \ldots, (a_n, b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^{k} a_i \le \sum_{i=1}^{k} \min(b_i, k-1) + \sum_{i=k+1}^{n} \min(b_i, k) \text{ holds for every } 1 \le k \le n.$ 

For simple polygon, when points are all integer, we have  $A = \#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$ .

Möbius inversion formula

$$\begin{array}{ll} - & f(n) = \sum_{d \mid n} g(d) \Leftrightarrow g(n) = \sum_{d \mid n} \mu(d) f(\frac{n}{d}) \\ - & f(n) = \sum_{n \mid d} g(d) \Leftrightarrow g(n) = \sum_{n \mid d} \mu(\frac{d}{n}) f(d) \end{array}$$

- Spherical cap
  - A portion of a sphere cut off by a plane.
  - -r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 h^2)$
  - Area =  $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- Lagrange multiplier
  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n) = 0$ .
  - Lagrangian function  $\mathcal{L}(x_1, \dots, x_n, \lambda_1, \dots, \lambda_k) = f(x_1, \dots, x_n)$  $\sum_{i=1}^{k} \lambda_i g_i(x_1, ..., x_n).$
  - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skewlines
  - $\operatorname{Line} 1: \boldsymbol{v}_1 = \boldsymbol{p}_1 + t_1 \boldsymbol{d}_1$  $- \operatorname{Line} 2: \boldsymbol{v}_2 = \boldsymbol{p}_2 + t_2 \boldsymbol{d}_2$
  - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
  - $-\mathbf{n}_1 = \mathbf{d}_1 \times \mathbf{n}$  $- \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$

  - $-\begin{array}{l} c_1 = p_1 + rac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1 \ c_2 = p_2 + rac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2 \end{array}$
- Derivatives/Integrals

Derivatives/Integrals Integration by parts: 
$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$
 
$$\left|\frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}}\right| \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}} \left|\frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2}\right|$$
 
$$\int_a^b \tan x = 1 + \tan^2 x \int_a^b \tan x = -\frac{\ln|\cos x|}{a}$$
 
$$\int_a^b e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \int_a^b xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$
 
$$\int_a^b \sqrt{a^2 + x^2} = \frac{1}{2}\left(x\sqrt{a^2 + x^2} + a^2 \sinh(x/a)\right)$$

$$(x,y,z) = (r\sin\theta\cos\phi, r\sin\theta\sin\phi, r\cos\theta)$$

$$(r,\theta,\phi) = (\sqrt{x^2 + y^2 + z^2}, a\cos(z/\sqrt{x^2 + y^2 + z^2}), a\tan(y,x))$$

Rotation Matrix

$$M(\theta) = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix}, R_x(\theta_x) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \cos\theta_x & -\sin\theta_x \\ 0 & \sin\theta & \cos\theta \end{bmatrix}$$

#### 6.16Estimation

n |2345678920304050100 p(n) = 23571115223062756044e42e52e8

n |1001e31e6 1e9 1e12 1e15 1e18 d(i) 12 32 240 1344 6720 26880 103680

 $n \mid 123456789$ 10 11 12 13 14 15  $\binom{2n}{n}$  2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8  $n \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid 10 \mid 11 \mid 12 \mid 13$  $B_n | 2515522038774140211471159757e54e63e7$ 

#### 6.17 Euclidean Algorithms

- $m = |\frac{an+b}{a}|$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c,b \bmod c,c,n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c,c - b - 1,a,m - 1) \\ - h(c,c - b - 1,a,m - 1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ &+ \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ &+ h(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ &- 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### General Purpose Numbers

Bernoulli numbers 
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 $S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$   $S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$   $x^n = \sum_{i=0}^{n} S(n,i)(x)_i$ • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. kj:ss.t.  $\pi(j) > \pi(j+1), k+1j$ :ss.t.  $\pi(j) \ge j$ , kj:ss.t.  $\pi(j) > j$ .

E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$ 

#### 6.19 Tips for Generating Functions

- Ordinary Generating Function  $A(x) = \sum_{i>0} a_i x^i$ 
  - $-A(rx) \Rightarrow r^n a_n$
  - $-A(x)+B(x) \Rightarrow a_n+b_n$

  - $-A(x)B(x)\Rightarrow \sum_{i=0}^{n}a_{i}b_{n-i}$   $A(x)^{k}\Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n}^{n}a_{i_{1}}a_{i_{2}}\dots a_{i_{k}}$

  - $-\stackrel{A(x)}{\xrightarrow{1-x}} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function  $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$ 
  - $-A(x)+B(x) \Rightarrow a_n+b_n$

  - $-A^{(k)}(x) \Rightarrow a_{n+k}$   $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
  - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $-xA(x) \Rightarrow na_n$
- Special Generating Function
  - $(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$
  - $-\frac{1}{(1-x)^n} = \sum_{i>0} {i \choose n-1} x^i$

## Polynomial

#### 7.1 Fast Fourier Transform [ec5a4e]

```
template<int MAXN>
struct FFT {
  \underline{using} \ val\_t = \underline{complex} < \underline{double} >;
  const double PI = acos(-1);
  val\_t\ w[M\!A\!X\!N]\,;
     for (int i = 0; i < MAXN; ++i) {
   double arg = 2 * PI * i / MAXN;
        w[i] = val_t(cos(arg), sin(arg));
  void bitrev(val_t *a, int n); // see NTT
        (val\_t *a, int n, bool inv = false); // see NIT;
   // remember to replace LL with val_t
};
```

#### Number Theory Transform\* [8fda91]

```
//(2^16)+1, 65537, 3
\frac{7}{7}*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881,
//51*(2^25)+1, 1711276033, 29
template<int MAXN, 11 P, 11 RT> //MAXN must be 2<sup>k</sup>
struct NTT {
  11 w [MAXN];
   ll mpow(ll a, ll n);
   11 \min(11 \ a) \{ return \operatorname{mpow}(a, P - 2); \}
  NTT() {
     11 \text{ dw} = \text{mpow}(RT, (P - 1) / MAXN);
     w[0] = 1;
     for (int
          i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw % P;
  void bitrev(ll *a, int n) {
     int i = 0;
     for (int j = 1; j < n - 1; ++j) {
for (int k = n >> 1; (i \hat{} = k) < k; k >>= 1);
if (j < i) swap(a[i], a[j]);
  void operator()(
        ll *a, int n, bool inv = false) { //0 \le a[i] \le P
     bitrev(a, n);

for (int L = 2; L <= n; L <<= 1) {

  int dx = MAXN / L, dl = L >> 1;
        for (int i = 0; i < n; i += L) {
          for (int
             j = i, x = 0; j < i + dl; ++j, x += dx) {

ll tmp = a[j + dl] * w[x] % P;
                   + dl = a[j] - tmp < 0 a[j + dl] += P;
             if ((a[j] += tmp) >= P) a[j] -= P;
          }
       }
     if (inv) {
       reverse (a + 1, a + n);
       ll invn = minv(n);
       for (int
              i = 0; i < n; ++i) a[i] = a[i] * invn % P;
  }
};
```

#### 7.3 Fast Walsh Transform\* [36c9f5]

```
/* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
or: (y += x * op), and: (x += y * op)

xor: (x, y = (x + y) * op, (x - y) * op)

invop: or, and, xor = -1, -1, 1/2 */

void fwt(int *a, int n, int op) { //or}

for (int L = 2; L <= n; L <<= 1)

for (int i = 0; i < n; i += L)

for (int j = i; j < i + (L >> 1); ++j)

a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
        N[1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
```

## 7.4 Polynomial Operation [37b8c7]

```
#define
using vector<ll>>::vector;
  static NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
Poly(const Poly &p, int m) : vector<ll>(m) {
     copy_n(p.data(), min(p.n(), m), data());
  Poly& irev()
  { return reverse(data(), data() + n()), *this; } Poly& isz(int m) { return resize(m), *this; } Poly& iadd(const Poly &rhs) { // n() == rhs.n()
     fi(0, n()) if
     (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
  Poly& imul(ll k) {
     fi(0, n()) (*this)[i] = (*this)[i] * k % P;
return *this;
  Poly Mul(const Poly &rhs) const {
     int m = 1;
     while (m < n() + rhs.n() - 1) m <<= 1;
     Poly X(*this, m), Y(rhs, m);
ntt(X.data(), m), ntt(Y.data(), m);
fi(0, m) X[i] = X[i] * Y[i] % P;
     ntt(X.data(), m, true);
return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
     if (n() == 1) return {ntt.minv((*this)[0])};
     while (m < n() * 2) m <<= 1;
     Poly Xi = Poly(*this, (n() + 1) / 2). Inv(). isz(m); Poly Y(*this, m);
     ntt(Xi.data(), m), ntt(Y.data(), m);
     fi(0, m) \{ Xi[i] *= (2 - Xi[i] * Y[i]) \% P;
        if ((Xi[i] %= P) < 0) Xi[i] += P;
     ntt(Xi.data(), m, true);
     return Xi.isz(n());
  Poly Sqrt()
         const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
     if (n()
          = 1) return {QuadraticResidue((*this)[0], P)};
     Poly
          X = Poly(*{\tt this} \;,\; (n() \;+\; 1) \;\;/\;\; 2) \,.\, Sqrt() \,.\, isz(n()) \,;
     return
           X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
  pair < Poly , Poly > DivMod
     (const Poly &rhs) const \{ // (rhs.)back() != 0 \} if (n() < rhs.n()) return \{ \{ 0 \}, *this \};
     const int m = n() - rhs.n() + 1;
     Poly X(rhs); X. irev(). isz(m);
     Poly Y(*this); Y. irev(). isz(m);
    Poly Q = Y.Mul(X.Inv()).isz(m).irev();

X = rhs.Mul(Q), Y = *this;

fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;

return {Q, Y.isz(max(1, rhs.n() - 1))};
  Poly Dx() const {
     Poly ret(n() - 1);
```

```
fi(0,
        ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
   return ret.isz(max(1, ret.n()));
Poly Sx() const {
   Poly ret(n() + 1);
   fi(0, n())
          ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
   return ret;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
    \begin{array}{lll} \textbf{return} & \textbf{Poly}(\grave{Y}.\,data() \, \stackrel{.}{+} \, \overset{.}{n()} \, \stackrel{.}{-} \, 1, \, Y.\, \overset{.}{data()} \, + Y.\,n()); \\ \end{array} 
vector<ll> _eval(const
       vector<ll> &x, const vector<Poly> &up) const {
   const int m = (int)x.size();
   if (!m) return {};
   vector < Poly > down (m * 2);
   // \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};
   // fi(2, m *
2) down[i] = down[i / 2].DivMod(up[i]).second;
   down[1] = Poly(up[1])
   \begin{array}{l} ..irev().isz(n()).Inv().irev().\_tmul(m, *this); \\ fi(2, m * 2) down[i] \\ = up[i \ \ 1].\_tmul(up[i].n() - 1, down[i / 2]); \end{array}
   vector < ll > y(m);
   fi(0, m) y[i] = down[m + i][0];
static vector<Poly> _treel(const vector<ll> &x) {
  const int m = (int)x.size();
   vector<Poly> up(m * 2);
   fi(0, m) up[m+i] = {(x[i] ? P - x[i] : 0), 1};
for (int i = m - 1; i
       > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
   return up;
}
vector
     <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
   auto up = \_tree1(x); return \_eval(x, up);
static Poly Interpolate (const vector
     <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
   const int m = (int)x.size();
  for (int i = m -
         1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
   return down[1];
Poly Ln() const { // (*this)[0] = 1, 1e5/170ms
   return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] = 0, 1e5/360ms
   if (n() = 1) return \{1\};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
   Poly Y = X.Ln(); Y[0] = P - 1;
   fi(0, n())
          if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;
   return X. Mul(Y). isz(n());
 V/M := P(P - 1). If k >= M, k := k \% M + M.
Poly Pow(ll k) const {
   int nz = 0;
  while (nz < n() && !(*this)[nz]) ++nz;

if (nz * min(k, (l1)n()) >= n()) return Poly(n());

if (!k) return Poly(Poly {1}, n());

Poly X(data() + nz, data() + nz + n() - nz * k);

const 1l c = ntt.mpow(X[0], k % (P - 1));
   return X.Ln().imul
         (k % P).Exp().imul(c).irev().isz(n()).irev();
static 11
      LinearRecursion(const vector<ll> &a, const vector
   const int k = (int)a.size();
const int k = (int)a.size();
  \begin{array}{l} assert \, ((\,int\,)\,coef.\,size\,() \,\, \Longrightarrow \,\, k \,+\, 1)\,; \\ Poly \,\, C(k \,+\, 1)\,, \,\, W(Poly \,\,\{1\}, \,\, k)\,, \,\, M \,=\, \{\,0\,,\,\, 1\,\}; \\ fi\,(\,1\,,\,\, k \,+\, 1)\,\, C[k \,-\, i\,] \,\, = \,coef\,[\,i\,] \,\,\,? \,\, P \,\, -\,\,coef\,[\,i\,] \,\,: \,\, 0; \end{array}
  C[k] = 1;
   while (n)
      if (n \% 2) W = W. Mul(M). DivMod(C). second;
```

```
n /= 2, M = M.Mul(M).DivMod(C).second;
}
ll ret = 0;
fi(0, k) ret = (ret + W[i] * a[i]) % P;
return ret;
}
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template ⇒ decltype(Poly_t::ntt) Poly_t::ntt = {};
```

#### 7.5 Value Polynomial [fad6e7]

```
struct Poly {
 mint base; // f(x) = poly[x - base]
 vector<mint> poly;
 Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
 mint get_val(const mint &x) {
   if (x \ge base \&\& x < base + SZ(poly))
    return poly[x - base];
   mint rt = 0;
   vector < mint > lmul(SZ(poly), 1), rmul(SZ(poly), 1);
   return rt;
 return:
   mint nw = get_val(base + SZ(poly));
   poly.pb(nw);
   for (int i = 1; i < SZ(poly); ++i)
     poly[i] += poly[i - 1];
};
```

#### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0\pmod{x^{2^k}}$ , then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

# 8 Geometry

#### 8.1 Basic [b068f0]

```
bool same
         (double a, double b) \{ return abs(a - b) < eps; \}
struct P {
    \quad \quad \text{double } x\,,\ y\,;
    P() : x(0), y(0) \{ \}
   P(): x(0), y(0) {}
P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }
double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
double obe() { return by by to (x y y ); }
    double abs() { return hypot(x, y); }
   P unit() { return *this / abs(); } P rot(double o) {
         double c = cos(o), s = sin(o);
return P(c * x - s * y, s * x + c * y);
    double angle() { return atan2(y, x); }
};
struct L {
    // ax + by + c = 0
    double a, b, c, o;
    P\ pa\,,\ pb\,;
   L(): a(0), b(0), c(0), o(0), pa(), pb() {}
L(P pa, P pb): a(pa.y - pb.y), b(pb.x - pa.x
), c(pa ^ pb), o(atan2(-a, b)), pa(pa), pb(pb) {}
    P project(P p) { return pa + (pb - pa).unit () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
```

```
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs() * (pb - pa).abs()); }
\max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
    if (\max(p1.y, p2.y) < \min(p3.y, p4.y) | |
           \max(p3.y, p4.y) < \min(p1.y, p2.y)) return false;
   (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
       (L x, L y) \{ return same(x.a * y.b, x.b * y.a); \}
       (L x, L y) { return P(-x.b * y.c + x.c * y.b, x .a * y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
       KD Tree [36d550]
\begin{array}{l} {\bf namespace} \;\; kdt \;\; \{\\ {\bf int} \;\; {\bf root} \;, \;\; lc \; [maxn] \;, \end{array}
         rc\left[ maxn\right] ,\ xl\left[ maxn\right] ,\ xr\left[ maxn\right] ,\ yl\left[ maxn\right] ,\ yr\left[ maxn\right] ;
point p[maxn];
int build(int 1, int r, int dep = 0) {
   if (l = r) return -1;
function<br/>
sol(const) point &, const point
           &>> f = [dep](const point &a, const point &b) {
       if (dep \& 1) return a.x < b.x;
       else return a.y < b.y;
    };
   int m = (l + r) >> 1;
   nth\_element\,(\,p\,+\,l\;,\;\;p\,+\,m,\;\;p\,+\,r\;,\;\;f\,)\;;
   xl[m] = xr[m] = p[m].x;

yl[m] = yr[m] = p[m].y;
   lc[m] = build(1, m, dep + 1);
   if (~lc[m]) {
       xl[m] = min(xl[m], xl[lc[m]]);
       \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \end{array} 
      yr[m] = max(yr[m], yr[lc[m]]);
   rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
       xl[m] = min(xl[m], xl[rc[m]]);
       \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{rc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{rc}\left[m\right]\right]\right); \\ \end{array} 
      yr[m] = max(yr[m], yr[rc[m]]);
bool bound(const point &q, int o, long long d) {
   double ds = sqrt(d + 1.0);
    if (q.x < xl[o] - ds | | q.x > xr[o] + ds | |
          q.y <
                 yl[o] - ds \mid \mid q.y > yr[o] + ds) return false;
   return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
void dfs (
      const point &q, long long &d, int o, int dep = 0) {
    if (!bound(q, o, d)) return;
   long long cd = dist(p[o], q);
if (cd != 0) d = min(d, cd);
    if ((dep & 1)
       && q.x < p[o].x || !(dep & 1) && q.y < p[o].y) { if (\simlc[o]) dfs(q, d, lc[o], dep + 1);
       if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
   } else {
       if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
void init(const vector<point> &v) {
   for (int i = 0; i < v.size(); ++i) p[i] = v[i];
   root = build(0, v.size());
long long nearest (const point &q) {
  long long res = 1e18;
```

```
dfs(q, res, root);
  return res;
8.3
       Sector Area [ec8913]
// calc area of sector which include a, b
double SectorArea(Pa, Pb, double r) {
  double o = atan2(a.y, a.x) - atan2(b.y, b.x);
  while (o \le 0) o += 2 * pi;
while (o \ge 2 * pi) o -= 2 * pi;
o = min(o, 2 * pi - o);
return r * r * o / 2;
8.4 Half Plane Intersection [0954c1]
```

```
bool jizz (L 11, L 12, L 13) {
  P p=Intersect (12, 13);
   return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;
bool cmp(const L &a, const L &b){
   return same(
        a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
   sort(ls.begin(),ls.end(),cmp);
   vector < L > pls(1, ls[0]);
   for (int i=0; i<(int) ls. size ();++i) if (!
        same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
   deque < int > dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c
     ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
   for(int i=2;i<(int) pls.size();++i){
  meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();</pre>
     meow(i,dq[0],dq[1])dq.pop\_front();
     dq.push_back(i);
  meow (dq
        .front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop\_front();
   if (dq.size()<3u)return vector
        <P>(); // no solution or solution is not a convex
   vector<P> rt;
   \begin{array}{ll} \text{for} (\, \text{int} \ i \! = \! 0; i \! < \! (\, \text{int}\,) \, dq \, . \, size \, (\,); \! + \! + \, i \,) \, rt \, . \, push\_back \end{array}
        (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
   return rt;
}
```

#### Rotating Sweep Line [b9fa8d]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
  int n=int(ps.size());
  vector < int > id(n), pos(n);
  vector < pair < int, int > line(n*(n-1)/2);
  int m=-1;
  for(int i=0;i< n;++i)for
       (int j=i+1; j< n; ++j) line[++m] = make\_pair(i,j); ++m;
  sort(line.begin(),line.end(),[&](const
        pair<int, int> &a, const pair<int, int> &b)->bool{
      if (ps
           [a.first].first=ps[a.second].first)return 0;
      if (ps
           [b.first].first=ps[b.second].first)return 1;
      return (double
           )(ps[a.first].second-ps[a.second].second)/(ps
           [a.first].first-ps[a.second].first) < (double
           )(ps[b.first].second-ps[b.second].second
)/(ps[b.first].first-ps[b.second].first);
      });
  for (int i=0; i< n; ++i) id [i]=i;
  sort(id.begin(),id.end(),[\&](const
        int &a, const int &b) { return ps[a] < ps[b]; });
  for (int i=0; i< n; ++i) pos [id[i]]=i;
  for (int i=0; i \triangleleft m; ++i)
    auto l=line[i];
    // meow
    tie (pos[l.first], pos[l.second],
         id [pos[l.first]], id [pos[l.second]])=make_tuple
         (pos[l.second],pos[l.first],l.second,l.first);
  }
```

#### 8.6 Triangle Center [33473a]

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
   Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)) / (\sin(a1) * \cos(a2) - \sin(a2) * \cos(a1));
return Point(ax + r1 * \cos(a1), ay + r1 * \sin(a1));
}
Point\ TriangleMassCenter(Point\ a,\ Point\ b,\ Point\ c)\ \{
  return (a + b + c) / 3.0;
Point\ TriangleOrthoCenter(Point\ a,\ Point\ b,\ Point\ c)\ \{
  return TriangleMassCenter(a, b
       , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
   double la = len(b - c);
   double lb = len(a - c);
  double lc = len(a - b);
  res.x = (
       la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
   res.y = (
la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

#### 8.7 Polygon Center [728c3a]

```
Point BaryCenter(vector<Point> &p, int n) {
  Point res(0, 0);
  double s = 0.0, t;
for (int i = 1; i < p.size() - 1; i++) {
    t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
    res.x += (p[0].x + p[i].x + p[i + 1].x)
    res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
  res.x /= (3 * s);
res.y /= (3 * s);
  return res;
```

#### 8.8 Maximum Triangle [55b8cb]

```
{\color{red} 	ext{double}} ConvexHullMaxTriangleArea
        (Point p[], int res[], int chnum) {
    double area = 0, tmp;
   res[chnum] = res[0];
for (int i = 0, j = 1, k = 2; i < chnum; i++) {
  while (fabs(Cross(p[</pre>
                \begin{array}{l} res \, [\,j\,]] \, \, - \, p \, [\, res \, [\,i\,]] \, , \, \, p \, [\, res \, [\,(\,k\,+\,1)\,\,\% \,\, chnum\,]] \, \, - \, \\ p \, [\, res \, [\,i\,]] \, , \, ) \, > \, fabs \, (Cross \, (p \, [\, res \, [\,j\,]] \, \, - \, p \, [\, res \, [\,i\,]] \, \, , \\ \end{array} 
                 p[res[k]] - p[res[i]]))) k = (k + 1) % chnum;
       tmp = fabs(Cross(
               p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
        if (tmp > area) area = tmp;
        while (fabs(Cross(p[
               res[(j+1)\% chnum]] - p[res[i]], p[res[k]]
               \begin{array}{l} p[\operatorname{res}[i]]) > \operatorname{fabs}(\operatorname{Cross}(p[\operatorname{res}[j]] - p[\operatorname{res}[i]], \\ p[\operatorname{res}[k]] - p[\operatorname{res}[i]]))) \ j = (j+1) \% \ \text{chnum}; \end{array}
       tmp = fabs(Cross(
               p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]);
        if (tmp > area) area = tmp;
   return area / 2;
```

#### 8.9 Point in Polygon [88cf80]

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) \% ps. size();
    L l(ps[a], ps[b]);
    P q = 1.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
    if (same(ps[
        a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
```

(ps[i], ps[j], r) \* (o>= 0 ? 1 : -1);

```
if \ (ps[a].y > ps[b].y) \ swap(a, \ b);\\
     if (ps[a].y <= p.y && p.y <
                                                                             return abs(ans);
          ps[a] \cdot y \leftarrow p \cdot y \text{ sace } p \cdot y

ps[b] \cdot y \text{ sace } p \cdot x \leftarrow ps[a] \cdot x + (ps[b] \cdot x - ps[a] \cdot x

) / (ps[b] \cdot y - ps[a] \cdot y) * (p \cdot y - ps[a] \cdot y)) ++c;
                                                                          8.11
                                                                                    Tangent of Circles and Points to Circle
  return (c & 1) * 2;
                                                                          vector<L> tangent(C a, C b) {
                                                                          #define Pij \
8.10 Circle [b6844a]
                                                                            P \ i \ = \ (b.c \ - \ a.c).unit() \ * \ a.r \, , \ j \ = \ P(\,i\,.y\,, \ -i\,.x\,)\,; \setminus
                                                                             z.emplace\_back(a.c + i, a.c + i + j);
struct C {
 Р с;
                                                                          #define deo(I,J)
                                                                             double d = (a
  double r:
  C(P \ c = P(0, \ 0), \ double \ r = 0) : c(c), r(r) \ \{\}
                                                                                 .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
                                                                             P i =
                                                                            (b.c - a.c).unit(), j = i.rot(o), k = i.rot(-o);\z.emplace_back(a.c + j * a.r, b.c J j * b.r);\z.emplace_back(a.c + k * a.r, b.c J k * b.r); if (a.r < b.r) swap(a, b);
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
                                                                             vector < L > z;
                                                                             if ((a.c - b.c).abs() + b.r < a.r) return z;
  if (same(a.r + b.r,
                                                                             else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
        d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
                                                                             else {
                                                                               deo(-,+);
     double o = acos
                                                                               if (same(d, a.r + b.r)) \{ Pij; \}
else if (d > a.r + b.r) \{ deo(+,-); \}
         ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
    P i = (b.c - a.c).unit();
p.push_back(a.c + i.rot(o) * a.r);
    p.push_back(a.c + i.rot(-o) * a.r);
                                                                             return z:
                                                                          }
  return p;
                                                                           vector <L> tangent (C c, P p) {
double IntersectArea(C a, C b) {
                                                                             vector<L> z;
                                                                             \frac{\text{double } d = (p - c.c).abs();}{}
  if (a.r > b.r) swap(a, b);
                                                                             if (same(d, c.r)) {
  double d = (a.c - b.c).abs();
  if (d >= a.r + b.r - eps) return 0;
                                                                               P i = (p - c.c).rot(pi / 2);
                                                                               z.emplace\_back(p, p + i);
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
                                                                             else if (d > c.r) {
  double p = acos
                                                                               double o = acos(c.r / d);
       ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
  double q = acos
                                                                               P i = (p - c.c).unit
                                                                                     (), j = i.rot(o) * c.r, k = i.rot(-o) * c.r;
       ((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
  z.emplace\_back(c.c + j, p);
                                                                               z.emplace\_back(c.c + k, p);
// remove second
      level if to get points for line (defalut: segment)
                                                                             return z;
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
                                                                                  Area of Union of Circles [0590f1]
                                                                          8.12
    double C = sq(a.x - o.x 
                                                                          vector<pair<double, double>>> CoverSegment(C &a, C &b) {
       ) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
                                                                             double d = (a.c - b.c).abs();
                                                                             vector<pair<double, double>>> res;
  vector<P> t;
  if (d >= -eps) {
                                                                              \begin{array}{ll} if \; (same(a.r + b.r, \; d)) \; ; \\ else \; if \; (d <= \; abs(a.r - b.r) + eps) \; \{ \\ if \; (a.r < b.r) \; res.emplace\_back(0, \; 2 \; * \; pi); \\ \end{array} 
     \mathbf{d} = \max(0., \mathbf{d});
     double i = (-B - sqrt(d)) / (2 * A);
double j = (-B + sqrt(d)) / (2 * A);
                                                                             } else if (d < abs(a.r + b.r) - eps) {
     if (i - 1.0 <= eps && i >=
                                                                               -eps) t.emplace_back(a.x + i * x, a.y + i * y);
     if (j - 1.0 <= eps && j >=
          -eps) t.emplace_back(a.x + j * x, a.y + j * y);
                                                                               double l = z - o, r = z + o;
                                                                               if (1 < 0) 1 += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
  return t;
                                                                               if (l > r) res.emplace_back
    (l, 2 * pi), res.emplace_back(0, r);
else res.emplace_back(l, r);
// calc area
      intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
                                                                             return res;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) -
                                                                          double CircleUnionArea
     if (inb) return abs(a ^ b) / 2;
                                                                                (\textit{vector} <\!\!\!\!\!\!<\!\!\!\!\!<\!\!\!c)\ \{\ /\!/\ \textit{circle should be identical}
     return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
                                                                             int n = c.size();
                                                                             double a = 0, w;
                                                                             for (int i = 0; w = 0, i < n; ++i) { vector<pair<double, double>> s = {{2 * pi, 9}}, z; for (int j = 0; j < n; ++j) if (i != j) {
  if (inb) return
  SectorArea(p[0], a, r) + abs(p[0] \hat{} b) / 2;
if (p.size() = 2u) return SectorArea(a, p[0], r)
+ SectorArea(p[1], b, r) + abs(p[0] \hat{} p[1])
                                                         p[1]) / 2;
                                                                                  z = CoverSegment(c[i], c[j]);
  else return SectorArea(a, b, r);
                                                                                  for (auto &e : z) s.push_back(e);
 // for any triangle
                                                                               double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
                                                                                for (auto &e : s) {
     int j = (i + 1) \% 3;
                                                                                  if (e.first > w) a += F(e.first) - F(w);
     double o = atan2
                                                                                  w = max(w, e.second);
     \begin{array}{lll} (ps\,[\,i\,].\,y,\ ps\,[\,i\,].\,x) \ \ -\ atan2\,(ps\,[\,j\,].\,y,\ ps\,[\,j\,].\,x)\,;\\ if\ (o>=\ pi)\ o=\ o-\ 2\ *\ pi\,;\\ if\ (o<=\ -pi)\ o=\ o+\ 2\ *\ pi\,; \end{array}
                                                                             return a * 0.5;
     ans += AreaOfCircleTriangle
```

#### 8.13 Minimun Distance of 2 Polygons [e9c988]

```
// p, q is convex double TwoConvexHullMinDist
   for (i =
        0; i < n; ++i) if(P[i].y < P[YMinP].y) YMinP = i;
   for (i =
         0\,;\;\; i\,<\,m;\; +\!\!+\!\!i\,) \quad \text{if}\; (Q[\,i\,]\,.\,y\,>\,Q[YM\!x\!Q]\,.\,y) \;\; YM\!x\!Q\,=\,i\;;
  P[n] = P[0], Q[m] = Q[0];
for (int i = 0; i < n; ++i) {
     while (tmp = Cross(
           \begin{array}{l} \text{Q[YMaxQ} + 1] & \text{-P[YMinP} + 1], \text{P[YMinP]} & \text{-P[YMinP} \\ + 1]) > \text{Cross}(\text{Q[YMaxQ]} & \text{-P[YMinP} + 1], \text{P[YMinP} \\ \end{array} 
           ] - P[YMinP + 1])) YMaxQ = (YMaxQ + 1) \% m;
     if (tmp < 0) ans = min(ans, PointToSegDist (P[YMinP], P[YMinP + 1], Q[YMaxQ]));
      \begin{array}{l} \textbf{else} \ \ ans = \min(ans \,, \ TwoSegMinDist(P[\\ YMinP] \,, \ P[YMinP \,+ \, 1] \,, \ Q[YMaxQ] \,, \ Q[YMaxQ \,+ \, 1])) \,; \end{array} 
     YMinP = (YMinP + 1) \% n;
  return ans;
}
8.14 2D Convex Hull [d97646]
bool operator < (const P &a, const P &b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator > (const P &a, const P &b) {
   return same(a.x, b.x) ? a.y > b.y : a.x > b.x;
#define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
   vector<P> p;
   sort(ps.begin(), ps.end(), [&] (Pa, Pb) { return
  same(a.x, b.x) ? a.y < b.y : a.x < b.x; }); for (int i = 0; i < ps. size(); ++i) {
     while (p.size() \ge 2 \& crx(p[p.size() -
           2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
     p.push_back(ps[i]);
   int t = p.size();
   for (int i = (int)ps.size() - 2; i >= 0; --i) {
     while (p.size() > t && crx(p[p.size() -
2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
     p.push_back(ps[i]);
  p.pop_back();
   return p;
int \ sgn(double
      x) { return same(x, 0) ? 0 : x > 0 ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n:
   vector<\!\!P\!\!> p,\ u,\ d;
  CH() {}
  CH(vector<P> ps) : p(ps) {
     n = ps.size();
     rotate (p. begin
           (), min_element(p.begin(), p.end()), p.end());
     auto t = max_element(p.begin(), p.end());
     d = \text{vector} \langle P \rangle (p. \text{begin}(), \text{next}(t));
     u = \text{vector} < P > (t, p.end()); u.push_back(p[0]);
   int find(vector < P > &v, P d) {
     int l = 0, r = v.size();
     while (1 + 5 < r) {

int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;

if (v[L] * d > v[R] * d) r = R;
        else l = L;
     for (int i = l +
            1; \ i < r\,; \ +\!\!\!+\!\! i\,) \ if \ (v\,[\,i\,] \ ^* \ d > v\,[\,x\,] \ ^* \ d) \ x = i\,;
     return x;
```

```
int findFarest(P v) {
       if (v.y > 0 \mid | v.y = 0 & v.x > 0) return
               ((int)d.size() - 1 + find(u, v)) \% p.size();
       return find(d, v);
      get(int l, int r, Pa, Pb) {
       int s = sgn(crx(a, b, p[1 \% n]));
       while (1 + 1 < r) {
          int m = (l + r) \gg 1;
          else r = m;
       return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
   vector<P> getLineIntersect(P a, P b) {
       int X = findFarest((b - a).rot(pi /
       int Y = findFarest((a - b).rot(pi / 2));
       if (X > Y) swap(X, Y);
       if (sgn
             \begin{array}{l} (crx(a,\ b,\ p[X]))\ *\ sgn(crx(a,\ b,\ p[Y])) < \ 0) \\ return\ \{get(X,\ Y,\ a,\ b),\ get(Y,\ X+n,\ a,\ b)\}; \end{array}
       return {}; // tangent case falls here
   void update_tangent(P q, int i, int &a, int &b) {
      if (sgn(crx(q, p[a], p[i])) > 0) a = i;
if (sgn(crx(q, p[b], p[i])) < 0) b = i;
   void bs(int 1, int r, Pq, int &a, int &b) {
       if (l == r) return
       update_tangent(q, 1 % n, a, b);
        \mbox{int } s = sgn(\,crx\,(q,\ p[\,l\,\,\%\,\,n]\,,\ p[(\,l\,+\,1)\,\,\%\,\,n]\,)\,)\,; 
       while (1 + 1 < r) {
          int m = (1 + r) >> 1;
          if (sgn(crx
                 (\,q\,,\,\,p\,[m\;\%\;\,n\,]\;,\,\,p\,[\,(m\,+\,\,1)\;\,\%\;\,n\,]\,)\,\,\Longrightarrow\,\,s\,)\;\,l\,\,=\,m;
          else r = m;
      update_tangent(q, r % n, a, b);
   int x = 1;
   for (int i = 1)
           + 1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
   return x;
int findFarest(P v) {
   if (v.y > 0 | | v.y = 0 & v.x > 0) return
            ((int)d.size() - 1 + find(u, v)) \% p.size();
   return find(d, v);
P get(int 1, int r, Pa, Pb) {
   int s = sgn(crx(a, b, p[1 \% n]));
    while (l + 1 < r) {
       int m = (1 + r) >> 1;
        \mbox{if } (sgn(crx(a,\ b,\ p[m\ \%\ n])) == s) \ l = m; \\
       else r = m;
   return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
vector<P> getIS(P a, P b) {
   \begin{array}{ll} & \text{int } X = \operatorname{findFarest}\left(\left(\begin{array}{cc} , & 1 \\ \end{array}\right).\operatorname{spin}\left(\left(\begin{array}{cc} pi & / & 2\right)\right);\\ & \text{int } Y = \operatorname{findFarest}\left(\left(\begin{array}{cc} a & - & b\right).\operatorname{spin}\left(\left(\begin{array}{cc} pi & / & 2\right)\right);\\ \end{array}\right. \end{array}
    \quad \text{if} \quad (X>Y) \ \operatorname{swap}(X,\ Y) \ ;
   \begin{array}{c} \text{if } (\operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[X])) \ * \ \operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[Y])) < \\ 0) \ \text{return } \{ \operatorname{get}(X,\ Y,\ a,\ b),\ \operatorname{get}(Y,\ X+n,\ a,\ b) \}; \end{array}
   return {};
void update_tangent(P q, int i, int &a, int &b) {
   \begin{array}{ll} \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{a}],\mbox{ }p[\mbox{i}])) > 0) \mbox{ }a = \mbox{i} \mbox{;} \\ \mbox{if } (\mbox{sgn}(\mbox{crx}(\mbox{q},\mbox{ }p[\mbox{b}],\mbox{ }p[\mbox{i}])) < 0) \mbox{ }b = \mbox{i} \mbox{;} \\ \end{array}
void bs(int l, int r, Pq, int &a, int &b) {
   if (l == r) return;
   update_tangent(q, l % n, a, b);
    \begin{array}{l} {\rm int} \ s = sgn(\,crx\,(q,\ p[\,l\ \%\ n]\,,\ p[\,(\,l\ +\ 1)\ \%\ n]\,)\,)\,; \end{array} 
   while (l + 1 < r) {
int m = (l + r) >> 1;
       if (sgn
              (\, crx \, (q,\ p \, [m \, \% \, \, n] \, , \ p \, [\, (m + \, 1) \, \, \% \, \, n \, ] \, ) \, = \!\!\!\! - \, s \, ) \ l \, = m;
       else r = m:
   update_tangent(q, r % n, a, b);
bool contain (P p) {
  if (p.x < d[0].x | | p.x > d.back().x) return 0;
```

```
auto it
       = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
   f (it->x = p.x) {
  if (it->y > p.y) return 0;
  else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
  it = lower_bound
       (u.\,begin\,()\;,\;u.\,end\,()\;,\;P(p.x,\;1e12)\;,\;greater<\!\!P\!\!>\!\!())\,;
  if (it->x = p.x) {
  if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
  return 1:
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
  a = b = 0;
  int i
       = lower\_bound(d.begin(), d.end(), p) - d.begin();
  bs(0, i, p, a, b);
bs(i, d.size(), p, a, b);
  i = lower_bound(
       u.begin(), u.end(), p, greater<P>()) - u.begin();
  bs((int
       )d. size() - 1, (int)d. size() - 1 + i, p, a, b);
  bs((int)d.size()
        -1 + i, (int)d.size() -1 + u.size(), p, a, b);
  return 1:
```

#### 8.15 3D Convex Hull [clae8f]

```
double
  absvol(const\ P\ a, const\ P\ b, const\ P\ c, const\ P\ d)\ \{return\ abs(((b-a)^(c-a))*(d-a))/6;
struct convex3D {
  static const int maxn=1010;
  struct T{
     int a,b,c;
     bool res;
     T()\{\}
     T(int a, int
           b, int c, bool res=1: a(a), b(b), c(c), res(res){}
  int n,m;
 P p[maxn];
T f[maxn*8];
  int id [maxn] [maxn];
  bool on (T &t, P &q) {
     return ((
          p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
  int g=id[a][b];
     if (f [g].res){
        if (on(f[g],p[q]))dfs(q,g);
          id [q][b]=id[a][q]=id[b][a]=m;
          f[m++]=T(b,a,q,1);
    }
  void dfs(int p, int i){
     f[i].res=0;
     \begin{array}{l} \text{meow}(p,f[\,i\,\,].\,b,f[\,i\,].\,a)\,;\\ \text{meow}(p,f[\,i\,\,].\,c,f[\,i\,\,].\,b)\,;\\ \text{meow}(p,f[\,i\,\,].\,a,f[\,i\,\,].\,c)\,;\\ \end{array} 
  void operator()(){
     if (n<4)return;
     _{i\,f\,([\,\&\,]\,()\,\{}
          for (int i=1;i< n;++i) if (abs
                (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
           }() || [&](){
           for(int | i=2;i<n;++i) if(abs((p[0]-p[i])
                 (p[1]-p[i]) > eps) return swap(p[2],p[i]),0;
          return 1;
}() || [&](){
           for (int i
                =3;i < n;++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
           return 1;
           }())return;
     for (int i=0; i<4;++i) {
       T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
        if(on(t,p[i]))swap(t.b,t.c);
```

```
id [t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
          f[m++]=t;
       for (int i=4; i< n; ++i) for
              (int j=0; j < m++j) if (f[j]. res & on(f[j], p[i])) {
           dfs(i,j);
          break:
       int mm=m; m=0;
       for (int i=0; i < mm++i) if (f[i].res) f[m++]=f[i];
    bool same(int i, int j){
       return !(absvol(p[f[i].a],p[f[i
              ].b], p[f[i].c], p[f[j].a])>eps || absvol(p[f[i].a], p[f[i].b], p[f[i].b])>eps || absvol
              (p[f[i].a], p[f[i].b], p[f[i].c], p[f[j].c])>eps);
    int faces(){
       int r=0;
       for (int i=0; i < m++i)
          int iden=1;
          \begin{array}{ll} \textbf{for} (\, \textbf{int} \ \ \textbf{j} \! = \! \textbf{0}; \textbf{j} \! < \! \textbf{i}; \! + \! + \! \textbf{j} \,) \, \textbf{if} (\, \textbf{same} (\, \textbf{i} \, , \, \textbf{j} \,) \,) \, \textbf{iden} \! = \! \textbf{0}; \end{array}
          r += iden;
       return r;
} tb;
```

#### 8.16 Minimum Enclosing Circle [7e5b31]

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
   double d = p0 ^ p1;
  return pt(x, y);
}
circle min_enclosing(vector<pt> &p) {
  random\_shuffle(p.begin(),\ p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {
     if (norm2(cent - p[i]) <= r) continue;</pre>
     \mathrm{cent}\,=\,p\,[\,\,i\,\,]\,;
     r = 0.0;
     for (int j = 0; j < i; ++j) {
  if (norm2(cent - p[j]) <= r) continue;</pre>
       cent = (p[i] + p[j])
       r = norm2(p[j] - cent);
       for (int k = 0; k < j; ++k) {
if (norm2(cent - p[k]) <= r) continue;
          cent = center(p[i], p[j], p[k]);
          r \, = \, norm2(\, p \, [\, k \, ] \  \, \text{- cent} \, ) \, ; \label{eq:reconstruction}
     }
  return circle(cent, sqrt(r));
```

#### 8.17 Closest Pair [7f292a]

```
double closest_pair(int l, int r) {
  // p should be sorted
       increasingly according to the x-coordinates.
  if (l = r) return 1e9;
  if (r - l = 1) return dist(p[l], p[r]);
  int m = (1 + r) >> 1;
  double d =
      min(closest\_pair(l, m), closest\_pair(m + 1, r));
  vector<int> vec;
  for (int i = m; i >= 1 &&
     fabs(p[m].x - p[i].x) < d; --i) vec.push_back(i);
     (int i = m + 1; i \le r \&\&
     fabs(p[m].x - p[i].x) < d; ++i) vec.push_back(i);
  sort(vec.begin(), vec.end()
  d = min(d, dist(p[vec[i]], p[vec[j]]));
   }
  return d;
```

#### 9 Else

#### 9.1 Cyclic Ternary Search\* [28a883]

```
/* bool pred(int a, int b); f(0) \sim f(n-1) \text{ is a cyclic-shift U-function} \\ \text{return idx s.t. pred}(x, \text{idx}) \text{ is false forall } x^*/\\ \text{int cyc\_tsearch(int } n, \text{ auto pred}) \ \{\\ \text{if } (n=1) \text{ return } 0;\\ \text{int } l=0, \text{ r=n; bool rv = pred}(1, 0);\\ \text{while } (r-1>1) \ \{\\ \text{int } m=(l+r) \ / \ 2;\\ \text{if } (\text{pred}(0,m) ? \text{ rv: pred}(m, (m+1) \% n)) \text{ } r=m;\\ \text{else } l=m;\\ \}\\ \text{return pred}(l, r \% n) ? l: r \% n;\\ \}
```

# 9.2 Mo's Algorithm (With modification) $_{[5dec12]}$

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  if (LBid != q.LBid) return LBid < q.LBid;
     if (RBid != q.RBid) return RBid < q.RBid;
    return T < b.T;
  }
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
  for (auto q: query) {
  while (T < q.T) addTime(L, R, ++T); // TODO
     while (T > q.T) subTime(L, R, T--); // TODO
    while (R < q.R) add(arr[++R]); // TODO while (L > q.L) add(arr[--L]); // TODO while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO
     // answer query
}
```

#### 9.3 Mo's Algorithm On Tree [4a7f74]

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) \operatorname{ord}[\operatorname{in}[u]] = \operatorname{ord}[\operatorname{out}[u]] = u
4) bitset MAXN> inset
struct Query {
   int L, R, LBid, lca;
Query(int u, int v) {
      int c = LCA(u, v);
      if (c = u \mid | c = v)
         q.lca \, = \, -1 \, , \, \, q.L \, = \, out \, [\, c \, \, \widehat{} \, \, u \, \, \widehat{} \, \, v \, ] \, , \, \, q.R \, = \, out \, [\, c \, ] \, ;
      else if (out[u] < in[v])

q.lca = c, q.L = out[u], q.R = in[v];
      else
      \begin{array}{l} q.\,lca \, = \, c \, , \ q.L \, = \, out \, [v] \, , \ q.R \, = \, in \, [u] \, ; \\ q.\,Lid \, = \, q.L \, \, / \  \, blk \, ; \end{array}
   bool operator < (const Query &q) const {
      if (LBid != q.LBid) return LBid < q.LBid;
      return R < q.R;
};
void flip(int x) {
   if (inset[x]) sub(arr[x]); // TODO
   else add(arr[x]); // TODO
      inset[x] = \sim inset[x];
void solve(vector<Query> query) {
   sort(ALL(query));
   int L = 0, R = 0;
```

```
22
              \begin{array}{ll} \mbox{while} & (R > q.R) & \mbox{flip} (\mbox{ord} [R--]) \,; \\ \mbox{while} & (L < q.L) & \mbox{flip} (\mbox{ord} [L++]) \,; \\ \end{array}
              if (~q.lca) add(arr[q.lca]);
              // answer query
              if (~q.lca) sub(arr[q.lca]);
       }
}
                   Additional Mo's Algorithm Trick
 9.4
  • Mo's Algorithm With Addition Only
           - Sort querys same as the normal Mo's algorithm.
           - For each query [l,r]:
            - If l/blk = r/blk, brute-force.
             - If l/blk \neq curL/blk, initialize curL := (l/blk + 1) \cdot blk, curR :=
                 curL-1
            - If r > curR, increase curR
             - decrease curL to fit l, and then undo after answering
  · Mo's Algorithm With Offline Second Time
           - Require: Changing answer \equiv adding f([l,r],r+1).
             - Require: f([l,r],r+1) = f([1,r],r+1) - f([1,l),r+1).
            - Part1: Answerall f([1,r],r+1) first.
                 Part2: Store curR \to R for curL (reduce the space to O(N)), and then
                  answer them\,by\,the\,second\,offline\,algorithm.
\mathbf{9.5}^{-} \begin{array}{l} \text{Note: You must do the above symmetrically for the left boundaries.}} \\ \mathbf{4.5} \\ \mathbf{4.5} \\ \mathbf{6.5} \\ \mathbf{7.5} \\ \mathbf{6.5} \\ \mathbf{7.5} \\ \mathbf{1.5} \\ \mathbf{
 ll hilbert(int n, int x, int y) {
        11 \text{ res} = 0;
        for (int s = n / 2; s; s >>= 1) {
             int rx = (x \& s) > 0;
             int ry = (y \& s) > 0;
res += s * 111 * s * ((3 * rx) ^ ry);
              if (ry = 0) {
                   if (rx = 1) x = s - 1 - x, y = s - 1 - y;
                   swap(x, y);
             }
       }
       return res;
 // n = 2^k 
 9.6 DynamicConvexTrick* [6a6f6d]
// only works for integer coordinates!! maintain max
 struct Line {
       mutable ll a, b, p;
       bool operator
                     <(const Line &rhs) const { return a < rhs.a; }</pre>
       bool operator <(ll x) const { return p < x; }
 struct DynamicHull : multiset<Line, less<>> {
        static const ll kInf = 1e18;
        ll Div(ll a,
                      ll b) { return a / b - ((a ^ b) < 0 && a % b); }
       bool isect (iterator x, iterator y) {
             if (y = end()) \{ x->p = kInf; return 0; \}
                          ->a = y->a) x->p = x->b > y->b ? kInf : -kInf;
              else x->p = Div(y->b - x->b, x->a - y->a);
             return x->p>= y->p;
       void addline(ll a, ll b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
  while (isect(y, z)) z = erase(z);
}
              if (x != begin
                           () & isect(--x, y) isect(x, y = erase(y));
              while ((y = x) != begin
                           () & (-x)-p = y-p isect(x, erase(y));
        11 query(ll x) {
  auto l = *lower_bound(x);
              return l.a * x + l.b;
};
                 All LCS* [ae68f0]
 void all_lcs(string s, string t) { // 0-base
        vector < int > h(SZ(t));
        iota(ALL(h), 0);
       for (int a = 0; a < SZ(s); ++a) {
              int v = -1;
              for (int c = 0; c < SZ(t); ++c)
                   if(s[a] = t[c] | | h[c] < v)
                         swap(h[c], v);
              // LCS(s[0, a], t[b, c]) = 
// c - b + 1 - sum([h[i] >= b] | i <= c)
```

// h[i] might become -1 !!

}

#### 9.8 AdaptiveSimpson\* [dc2085]

```
template<typename Func, typename d = double>
struct Simpson {
  \begin{array}{ll} \textbf{using} \ \ \textbf{pdd} \ = \ \textbf{pair} <\!\! \textbf{d} \,, \ \ \textbf{d} \!\! > ; \end{array}
  Func f:
  d \ eval(pdd \ l, \ pdd \ r, \ d \ fm, \ d \ eps) \ \{ \ pdd \ m((l.X + r.X) \ / \ 2, \ fm);
     d s = mix(1, r, fm).second;
     auto [flm, sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
d delta = sl + sr - s;
     if (abs(delta ) <= 15 * eps) return sl + sr + delta / 15;
     return eval(1, m, flm, eps / 2) +
        eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
     return eval
          (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
  d \text{ eval2}(d 1, d r, d \text{ eps}, \text{ int } k = 997)  { d h = (r - 1) / k, s = 0;
     for (int i = 0; i < k; ++i, l += h)
       s \leftarrow eval(l, l + h, eps / k);
     return s;
  }
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
```

#### 9.9 Simulated Annealing [b14262]

```
double factor = 1000000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 10000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.999995;
}
```

#### 9.10 Tree Hash\* [e57357]

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
        sum += shift(dfs(i, u));
    return sum;
}</pre>
```

#### 9.11 Binary Search On Fraction [951597]

```
struct Q {
  ll p, q;
  Q go(Q b, 11 d) \{ return \{p + b.p*d, q + b.q*d\}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
   pred(p/q) is true, and 0 \le p,q \le N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
assert(pred(hi));
   bool dir = 1, L = 1, H = 1;
   for (; L | | H; dir = !dir) {
      11 len = 0, step = 1;
      for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
         \begin{array}{l} \text{if } \left(Q \text{ mid} = \text{hi.go}(\text{lo}\,,\, \text{len} + \text{step})\,; \\ \text{mid.p} > N \mid \mid \text{mid.q} > N \mid \mid \text{dir} \, \widehat{} \, \text{pred}(\text{mid})\right) \end{array}
         else len += step;
      swap (\,lo\;,\;\; hi\;=\; hi\,.\,go (\,lo\;,\;\; len\,)\,)\;;
      (dir ? L : H) = !! len;
```

```
}
return dir ? hi : lo;
}
```

#### 9.12 Bitset LCS [a82d86]

```
\begin{array}{lll} & \text{cin} >> n >> m; \\ & \text{for (int } i = 1, \ x; \ i <= n; \ +\!\!+\!\!i) \\ & \text{cin} >> x, \ p[x]. set(i); \\ & \text{for (int } i = 1, \ x; \ i <= m; \ i\!\!+\!\!+\!\!) \ \{ \\ & \text{cin} >> x, \ (g = f) |= p[x]; \\ & \text{f.shiftLeftByOne()}, \ f. set(0); \\ & ((f = g - f) \ \hat{} = g) \&= g; \\ & \text{cout} << f. count() << '\n'; \end{array}
```

# 10 Python

#### 10.1 Misc