Contents

```
13
  6.12Simplex Algorithm .
                                                               13
                                        6.12.1Construction .
                                    6.13 chinese Remainder . .
                                                               14
                                    6.14Factorial without prime factor*....
6.15Discrete Log*....
  1.4 Black Magic . . .
  1.5 Pragma Optimization
  1.6 Bitset . . . . . . . .
                                    6.16Berlekamp Massey .
  Graph
                                    6.17Primes . . . . . . . .
  6.18Theorem . . . . . . .
                                                               14
  6.19Estimation . . . . . .
                                                               15
                                    6.20Euclidean Algorithms
                                                               15
  {\bf 2.5~MinimumMeanCycle*}
                                    6.21General Purpose
  2.6 Virtual Tree*
                                        Numbers . . . . . . .
                                                               15
  2.7 Maximum Clique Dyn*
                                    6.22Tips for Generating
                  Steiner
  2.8 Minimum
                                        Functions . . . . . .
                                                               15
      Tree* . . .
  7 Polynomial
  \begin{array}{cccc} \textbf{2.10Minimum} & \textbf{Clique} \\ \textbf{Cover*} & \dots & \dots & \dots \end{array}
                                    7.1 Fast Fourier Transform 15
                                        \begin{array}{ll} Number & Theory \\ Transform^* \ . \ . \ . \ . \ . \end{array}
                                    7.2 Number
  2.11NumberofMaximalClique*
                                    7.3 Fast Walsh Transform* 16
3 Data Structure
  3.1 Discrete Trick . . . . 3.2 BIT kth* . . . . . . . .
                                    7.4 Polynomial Operation 16
                                    7.5 Value Polynomial . .
                                                               17
  3.3 Interval Container* .
3.4 Leftist Tree . . . . .
                                    7.6 Newton's Method . .
                                                               17
  3.5 Heavy light Decom-
                                    Geometry
      position*
  5
                                    8.1 Basic . . . . . . . .
                                                               17
                                    8.2 KD Tree . . . . . .
                                                               18
  8.3 Sector Area . . . . .
                                                               18
                                    8.4 Half Plane Intersection 18
                                    8.5 Rotating Sweep Line
                                                               18
  3.9~\mathrm{KDTree} . . . . . .
                              6
                                    8.6 Triangle Center . . .
                                                               19
  Flow/Matching
4.1 Bipartite Matching*
                                    8.7 Polygon Center . . .
                                                               19
                                    8.8 Maximum Triangle .
                                                               19
  4.3 MincostMaxflow* . . . 4.4 Maxim
  4.2 Kuhn Munkres*
                                    8.9 Point in Polygon . . .
                                                               19
                                    8.10Circle . . . . . . . . . . . 8.11Tangent of Circles
  4.4 Maximum
                   Simple
      Graph Matching*
                              7
                                        and Points to Circle.
  4.5 Maximum Weight Matching* . . . . . .
                                    8.12Area of Union of Circles 20
                                    8.13Minimun Distance of
  4.6 SW-mincut.
                                        2 Polygons . . . . . .
  4.7 BoundedFlow*(Dinic*)
                                    8.142D Convex Hull . . .
                                                               20
  4.8 Gomory Hu tree*
  8.153D Convex Hull . . .
                                                               21
                                    8.16Minimum Enclosing
                             10
                                        Circle . .
                                    8.17Closest Pair . . . . .
  String
                             10
  10
  9
                                    Else
                             10
                                    9.1 Cyclic Ternary Search* 22
                             10
                                    9.2 Mo's
                                                       Algo-
  11
                                        rithm(With modifi-
                                        cation)
                                                               22
      tomatan* . . . . . .
                                    9.3 Mo's Algorithm On
  5.6 Smallest Rotation . .
                             11
                                                               22
  5.7 De Bruijn sequence*
                                        Tree . . . . . . . . . .
                                    9.4 Additional Mo's Al-
  5.8 Extended SAM* . . .
                                        gorithm Trick . . . .
  5.9 PalTree* . . . . . .
  5.10Main Lorentz . . . .
                                    9.5 Hilbert Curve . .
                                    9.6 DynamicConvexTrick*
 Math
                             12
                                    9.7 All LCS* . . . . . . .
                                                               23
  9.8 AdaptiveSimpson* .
                                                               23
                                    9.9 Simulated Annealing
                                                               23
                                    9.10Tree Hash* . . . . . .
                                                               23
  6.3 Floor Enumeration .
                                    9.11Binary Search On
Fraction . . . . . . .
  6.4 Mod Min . . . . . .
                             12
  6.5 Linear Mod Inverse .
                                    9.12Bitset LCS . . . . . .
  6.6 Linear Filter Mu . .
                             13
  6.7 Gaussian integer gcd
                             13
  6.8 GaussElimination . .
                                  10 Python
  6.9 Miller Rabin* . . . .
                                    10.1Misc . . . . . . . . .
```

6.10Simultaneous Equa-

Basic 1

1.1vimrc

```
"This file should be placed at ~/.vimrc"
se nu ai hls et ru ic is sc cul
se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
syntax on
hi cursorline cterm=none ctermbg=89
set bg=dark
inoremap {<CR> {<CR>}<Esc>ko<tab>
    region and then type : Hash to hash your selection."
"Useful for verifying that there aren't mistypes. ca Hash w !cpp -dD -P -fpreprocessed
      \| tr -d '[:space:]' \| md5sum \| cut -c-6
```

1.2 readchar [0754b0]

```
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
static char *p = buf, *end = buf;
  if (p == end) end = buf +
         fread_unlocked(buf, 1, bufsize, stdin), p = buf;
  return *p++;
}
1.3 BigIntIO [ea947e]
  _int128 read() {
       _{\text{int}128} \text{ x} = 0, \text{ f} = 1;
     char ch = getchar();
     while (ch < '0' | | ch > '9') {
    if (ch == '-') f = -1;
          ch = getchar();
     while (ch >= '0' && ch <= '9') {
 x = x * 10 + ch - '0';
          ch = getchar();
     \textcolor{return}{\texttt{return}} \ x \ * \ f;
void print(__int128 x) {
     if (x < 0) {
          putchar('-');
          x = -x;
     if (x > 9) print (x / 10);
putchar (x \% 10 + '0');
bool cmp(\underline{\text{int}128 x}, \underline{\text{int}128 y}) { return x > y; }
1.4 Black Magic [d566f1]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace _
                      _gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
      tree{<} int \;,\; null\_type \;,\; std::less{<} int{>},\; rb\_tree\_tag
       tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int , int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
(0) == 22); assert(*s.find_by_order(1) == 71);
   assert(s.order_of_key
       (22) = 0; assert(s.order_of_key(71) == 1);
   s.erase(22);
  assert(*s.find_by_order
(0) == 71); assert(s.order_of_key(71) == 0);
    / mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0:
```

1.5 Pragma Optimization [7b330a]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse, sse2, sse3, ssse3, sse4")
#pragma GCC target("popent, abm, mmx, avx, arch=skylake")
   \_builtin\_ia32\_ldmxcsr(\_\_builtin\_ia32\_stmxcsr() | 0 \times 8040)
```

Bitset [282252]

```
#include < bits / stdc++.h>
using namespace std;
int main () {
    bitset <4> bit;
```

ECC(int _n): n(_n), dft()

, ecnt(), necc(), low(n), dfn(n), bln(n), G(n) {} void add_edge(int u, int v) {}

```
bit.all(); // all bit is true, ret tru;
bit.any(); // any bit is true, ret true
bit.none(); // all bit is false, ret true
bit count();
                                                                                                                                        G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
                                                                                                                                    void solve() {
                                                                                                                                        is_bridge.resize(ecnt);
for (int i = 0; i < n; ++i)
         bit.count();
         bit.to_string('0', '1');//with parmeter
         bit.reset(); // set all to true
bit.set(); // set all to false
                                                                                                                                             if (!dfn[i]) dfs(i, -1);
         std::bitset <8> b3{0}, b4{42};
                                                                                                                               }; // ecc_id(i): bln[i]
         std::hash<std::bitset<8>> hash_fn8;
                                                                                                                               2.3 SCC* [22afe1]
         hash_fn8(b3); hash_fn8(b4);
}
                                                                                                                               struct SCC { // 0-base
  int n, dft, nscc;
  vector<int> low, dfn, bln, instack, stk;
\mathbf{2}
          Graph
2.1
           BCC Vertex* [ed8308]
                                                                                                                                    vector<vector<int>>> G;
                                                                                                                                    void dfs(int u) {
struct BCC { // 0-base
int n, dft, nbcc;
                                                                                                                                        low[u] = dfn[u] = ++dft;
                                                                                                                                        \begin{array}{l} instack \left[u\right] = 1, \ stk.pb(u); \\ for \ (int \ v \ : G[u]) \end{array}
     vector<int> low, dfn, bln, stk, is_ap, cir;
     vector<vector<int>>> G, bcc, nG;
                                                                                                                                             if (!dfn[v])
     void make_bcc(int u) {
                                                                                                                                         \begin{array}{l} dfs(v), \ low[u] = min(low[u], \ low[v]); \\ else \ if \ (instack[v] \&\& \ dfn[v] < dfn[u]) \\ low[u] = min(low[u], \ dfn[v]); \\ if \ (low[u] = dfn[u]) \end{array}
         bcc.emplace_back(1, u);
for (; stk.back() != u; stk.pop_back())
  bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
         stk.pop\_back(), bln[u] = nbcc++;
                                                                                                                                             for (; stk.back() != u; stk.pop_back())
                                                                                                                                                 bln[stk
     void dfs(int u, int f) {
                                                                                                                                                           .back()] = nscc, instack[stk.back()] = 0;
         int child = 0;
                                                                                                                                             instack[u] = 0, bln[u] = nscc++, stk.pop\_back();
         low[u] = dfn[u] = ++dft, stk.pb(u);
                                                                                                                                        }
         for (int v : G[u])
             if (!dfn[v]) {
                                                                                                                                    SCC(int _n): n(_n), dft(), nscc
                 dfs(v, u), ++child;

low[u] = min(low[u],
                                                                                                                                             (), low(n), dfn(n), bln(n), instack(n), G(n) {}
                                                              low[v]);
                                                                                                                                    void add_edge(int u, int v) {
                   \begin{array}{ll} & \text{if } (dfn [u] <= low [v]) \\ & \text{is\_ap}[u] = 1, \ bln [u] = nbcc; \\ & \text{make\_bcc}(v), \ bcc.back().pb(u); \\ \end{array} 
                                                                                                                                        G[u].pb(v);
                                                                                                                                    void solve() {
  for (int i = 0; i < n; ++i)</pre>
              else \ if \ (dfn[v] < dfn[u] & v != f)
                                                                                                                                             if (!dfn[i]) dfs(i);
        \begin{array}{l} low[u] = min(low[u], dfn[v]); \\ if \ (f = -1 \&\& child < 2) \ is\_ap[u] = 0; \\ if \ (f = -1 \&\& child = 0) \ make\_bcc(u); \end{array}
                                                                                                                               }; // scc_id(i): bln[i]
                                                                                                                                2.4 2SAT* [e839e5]
    BCC(int _n): n(_n), dft()
                                                                                                                               struct SAT { // 0-base
             nbcc(), low(n), dfn(n), bln(n), is_ap(n), G(n)  {}
    void add_edge(int u, int v) {
  G[u].pb(v), G[v].pb(u);
                                                                                                                                    int n;
                                                                                                                                    vector<bool> istrue;
                                                                                                                                    SCC scc;
     void solve() {
  for (int i = 0; i < n; ++i)</pre>
                                                                                                                                    SAT(int \underline{n}): n(\underline{n}), istrue(n + n), scc(n + n)  {}
                                                                                                                                    int rv(int a) {
                                                                                                                                        return a >= n? a - n : a + n;
              if (!dfn[i]) dfs(i, -1);
                                                                                                                                    void add_clause(int a, int b) {
    scc.add_edge(rv(a), b), scc.add_edge(rv(b), a);
     void block_cut_tree() {
         cir.resize(nbcc);
         for (int i = 0; i < n; ++i)
              if (is_ap[i])
                                                                                                                                    bool solve()
                b \ln [i] = nbcc++;
                                                                                                                                        scc.solve();
                                                                                                                                        for (int i = 0; i < n; ++i) {
  if (scc.bln[i] == scc.bln[i + n]) return false;</pre>
         cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
              for (int j : bcc[i])
                                                                                                                                             istrue[i] = scc.bln[i] < scc.bln[i + n];
    if (is_ap[j])
    nG[i].pb(bln[j]), nG[bln[j]].pb(i);
} // up to 2 * n - 2 nodes!! bln[i] for id
                                                                                                                                             istrue[i + n] = !istrue[i];
                                                                                                                                        return true:
                                                                                                                               };
2.2 Bridge* [Occada]
                                                                                                                                             MinimumMeanCycle* [4be648]
                                                                                                                                2.5
{\tt struct} \ ECC \ \{ \ // \ 0\text{-base}
    int n, dft, ecnt, necc;
vector<int> low, dfn, bln, is_bridge, stk;
                                                                                                                                ll road [N] [N]; // input here
                                                                                                                                struct MinimumMeanCycle {
                                                                                                                                    11 dp[N + 5][N], n;
     vector<vector<pii>>> G;
    void dfs(int u, int f) {
dfn[u] = low[u] = ++dft, stk.pb(u);
                                                                                                                                     pll_solve() {
                                                                                                                                        11 \ a = -1, \ b = -1, \ L = n + 1;
         for (auto [v, e] : G[u])
if (!dfn[v])
                                                                                                                                         for (int i = 2; i \le L; ++i)
                                                                                                                                             for (int k = 0; k < n; ++k)
         dfs(v, e), low[u] = min(low[u], low[v]);

else if (e!= f)

low[u] = min(low[u], dfn[v]);

if (low[u] == dfn[u]) {
                                                                                                                                                 for (int j = 0; j < n; ++j)
                                                                                                                                                     dp[i][j] =
                                                                                                                                        fighting in the second in
              if (f != -1) is_bridge[f] = 1;
for (; stk.back() != u; stk.pop_back())
                                                                                                                                             11 ta = 0, tb = 1;
                                                                                                                                            for (int j = 1; j < n; ++j)
  if (dp[j][i] < INF &&
    ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
    ta = dp[L][i] - dp[j][i], tb = L - j;</pre>
                 bln[stk.back()] = necc;
              bln[u] = necc++, stk.pop\_back();
        }
```

if (ta = 0) continue;

if (a = -1) | a * tb > ta * b) a = ta, b = tb;

```
if (a != -1) {
    ll g = __gcd(a, b);
    return pll(a / g, b / g);
}
    return pll(-1LL, -1LL);
}
void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
}
};</pre>
```

2.6 Virtual Tree* [80f7cb]

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
   if (top = -1) return st[++top] = u, void();
   int p = LCA(st[top], u);
  if (p = st[top]) return st[++top] = u, void(); while (top >= 1 && dep[st[top - 1]] >= dep[p])
   vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
     \overrightarrow{vG}[p] \cdot \overrightarrow{pb}(\overrightarrow{st}[\overrightarrow{top}]), --top, \overrightarrow{st}[++top] = p;
   st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top =
   \operatorname{sort}\left(ALL(v)\right.,
     [\&](int a, int b) \{ return dfn[a] < dfn[b]; \});
   for (int i : v) insert(i);
   while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
   // do something
   reset(v[0]);
```

2.7 Maximum Clique Dynst [4a6b3d]

```
struct MaxClique { // fast when N <= 100
  bitset \langle N \rangle G[N], cs[N];
int ans, sol[N], q, cur[N], d[N], n;
void init(int _n) {
    n\,=\,\underline{\phantom{a}}n\,;
     for (int i = 0; i < n; ++i) G[i].reset();
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  \begin{tabular}{ll} void & pre\_dfs(vector < int > \&r \,, & int \ l \,, & bitset < N > mask) & \{ \end{tabular} \label{table}
     if (l < 4) {
       for (int i : r) d[i] = (G[i] \& mask).count();
       sort (ALL(r)
             , [\&](int x, int y) \{ return d[x] > d[y]; \});
     vector < int > c(SZ(r));
     int lft = \max(\text{ans} - q + 1, 1), rgt = 1, tp = 0;
     cs[1].reset(), cs[2].reset();
     for (int p : r) {
       int k = 1:
       while ((cs[k] \& G[p]).any()) ++k;
       if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
       if (k < lft) r[tp++] = p;
     for (int k = lft; k \ll rgt; ++k)
       for (int p = cs[k]._Find_first
             ()\;;\;\;p<\;N;\;\;p=\;cs\left[\,k\,\right].\,\_Find\_next\left(\,p\,\right)\,)
          r[tp] = p, c[tp] = k, +tp;
     dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
       int>&r, vector<int>&c, int l, bitset<N> mask) {
     while (!r.empty()) {
       int p = r.back();
       r.pop\_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;
       \operatorname{cur}[q++] = p;
       vector<int> nr;
       for (int i : r) if (G[p][i]) nr.pb(i);
```

```
if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
    else if (q > ans) ans = q, copy_n(cur, q, sol);
    c.pop_back(), --q;
}
}
int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
}
};
```

2.8 Minimum Steiner Tree* [cbf811]

```
struct SteinerTree { // 0-base
  \begin{array}{l} int \ n, \ dst\left[N\right]\left[N\right], \ dp\left[1 << T\right]\left[N\right], \ tdst\left[N\right]; \\ int \ vcst\left[N\right]; \ // \ the \ cost \ of \ vertexs \end{array}
   void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {
       fill_n (dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
    x = \min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
     for (int k = 0; k < n; ++k)
       for (int i = 0; i < n; ++i)
          for (int j = 0; j < n; ++j)
            chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1; for (int i = 0; i <= full; ++i)
       fill_n (dp[i], n, INF);
     copy_n(vcst, n, dp[0]);
     for (int msk = 1; msk \le full; ++msk) {
       if (!(msk & (msk - 1))) {
          int who = 1
                        _{
m lg}({
m msk});
          for (int i = 0; i < n; ++i)
            dp [msk
                 ][i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; +++i)
          for (int sub = (
               msk - 1) & msk; sub; sub = (sub - 1) & msk)
            chmin (dp [msk] [i],
                dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {
          tdst[i] = INF;
          for (int j = 0; j < n; +++j)
            chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
; // O(V 3^T + V^2 2^T)
```

2.9 Dominator Tree* [e95beb]

```
struct dominator_tree { // 1-base
  vector < int > G[N] , rG[N];
  int n, pa[N] , dfn[N] , id[N] , Time;
  int semi[N] , idom[N] , best[N];
  vector < int > tree [N]; // dominator_tree
  void init (int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
        G[i].clear(), rG[i].clear();
  }
  void add_edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  }
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
        if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  }
  int find(int y, int x) {</pre>
```

```
if (y \le x) return y;
      int tmp = find (pa[y], x);
      \begin{array}{l} \text{if } (\text{semi} [\text{best}[y]] > \text{semi} [\text{best}[pa[y]]]) \\ \text{best}[y] = \text{best}[pa[y]]; \end{array}
      return pa[y] = tmp;
   void tarjan(int root) {
      Time = 0;
      for (int i = 1; i \le n; ++i) {
         dfn[i] = idom[i] = 0;
         tree[i].clear();
best[i] = semi[i] = i;
      dfs(root);
      for (int i = Time; i > 1; --i) {
         \begin{array}{ll} {\bf int} \  \, u \, = \, id \, [\, i \, ] \, ; \end{array}
         for (auto v : rG[u])
            if (v = dfn[v]) 
              find(v, i);
semi[i] = min(semi[i], semi[best[v]]);
         tree[semi[i]].pb(i);
         for (auto v : tree[pa[i]]) {
           find(v, pa[i]);
idom[v] =
              semi[best[v]] = pa[i] ? pa[i] : best[v];
         tree [pa[i]]. clear();
      for (int i = 2; i \leftarrow Time; ++i)
         if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
tree[id[idom[i]]].pb(id[i]);
  }
};
```

2.10 Minimum Clique Cover* [5951ca]

```
struct Clique_Cover { // 0-base, O(n2^n)
  int dp[1 << N];
  void init(int _n) {
     n = n, fill_n (dp, 1 << n, 0);
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
     E[u] = 1 << v, E[v] = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)

co[1 << i] = E[i] | (1 << i);

co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {
        int t = i & -i;

dp[i] = -dp[i ^ t];

co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)
        co[i] = (co[i] \& i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)
    sum += (dp[i] *= co[i]);</pre>
        if (sum) return ans;
     }
     return n;
  }
};
```

2.11 NumberofMaximalClique* [c163d7]

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all [N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
        for (int j = 1; j <= n; ++j) g[i][j] = 0;
  }
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  }
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
```

```
int u = some[d][0];
for (int i = 0; i < sn; ++i) {
    int v = some[d][i];
    if (g[u][v]) continue;
    int tsn = 0, tnn = 0;
    copy_n(all[d], an, all[d + 1]);
    all[d + 1][an] = v;
    for (int j = 0; j < sn; ++j)
        if (g[v][some[d][j]])
            some[d + 1][tsn++] = some[d][j];
        for (int j = 0; j < nn; ++j)
            if (g[v][none[d][j]])
            none[d + 1][tnn++] = none[d][j];
        dfs(d + 1, an + 1, tsn, tnn);
        some[d][i] = 0, none[d][nn++] = v;
    }
}
int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
}
};</pre>
```

3 Data Structure

3.1 Discrete Trick

3.3 Interval Container* [78516e]

}

```
remove intervals from a set of disjoint intervals.
 * Will merge the added interval with
       any overlapping intervals in the set when adding.
 st Intervals are [inclusive, exclusive). st_{/}
set<pii>>::
     iterator addInterval(set<pii>% is, int L, int R) {
  if (L == R) return is.end()
  \begin{array}{lll} \textbf{auto} & \textbf{it} = \textbf{is.lower\_bound}(\{L,\ R\})\,, \ \textbf{before} = \textbf{it}\,; \end{array}
  while (it != is.end() & it->X<= R) {
    R = \max(R, it ->Y);
     before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y>= L) {
    L = \min(L, it ->X);
    R = \max(R, it ->Y);
    is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii> is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  \begin{array}{ll} \textbf{auto} & r2 \ = \ it \text{->}Y; \end{array}
  if (it->X == L) is .erase(it);
  else (int\&)it ->Y = L;
  if (R != r2) is .emplace(R, r2);
```

3.4 Leftist Tree [bbd228]

```
struct node {
    ll v, data, sz, sum;
    node *1, *r;
    node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
```

```
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
if (a->data < b->data) swap(a, b);
  a \rightarrow r = merge(a \rightarrow r, b);
   if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
   return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
   delete tmp;
3.5 Heavy light Decomposition* [babe8a]
```

```
\begin{array}{lll} \textbf{struct} & \textbf{Heavy\_light\_Decomposition} ~\{~//~~ \textbf{1-base} \\ & \textbf{int} ~n, ~u \textbf{link} [N] ~, ~d \textbf{eep} [N] ~, ~m \textbf{xson} [N] ~, ~w [N] ~, ~p \textbf{a} [N] ~; \end{array}
   int t, pl[N], data[N], val[N]; // val: vertex data
   vector<int>G[N];
   void init(int _n) {
      n = _n;
for (int i = 1; i <= n; ++i)
         G[i].clear(), mxson[i] = 0;
   void add_edge(int a, int b) {
      G[a].pb(b), G[b].pb(a);
   void dfs(int u, int f, int d) {
  w[u] = 1, pa[u] = f, deep[u] = d++;
  for (int &i : G[u])
          if (i != f) {
            dfs(i, u, d), w[u] += w[i];
             if (w[mxson[u]] < w[i]) mxson[u] = i;
   }
    \begin{array}{l} \mbox{void } \mbox{cut(int } \mbox{u, int } \mbox{link)} \ \{ \\ \mbox{data[pl[u] = ++t] = val[u], ulink[u] = link;} \end{array} 
      if (!mxson[u]) return;
      cut(mxson[u], link);
for (int i : G[u])
          if (i != pa[u] & i != mxson[u])
            cut(i, i);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
      int ta = ulink[a], tb = ulink[b], res = 0;
while (ta != tb) {
          if (deep
             [ta] > deep[tb]) swap(ta, tb), swap(a, b);
query(pl[tb], pl[b])
         tb = ulink[b = pa[tb]]
      if (pl[a] > pl[b]) swap(a, b);
       // query(pl[a], pl[b])
};
```

Centroid Decomposition* [4eccaf] 3.6

```
struct Cent_Dec { // 1-base
  vector<pll> G[N];
  pll info [N]; // store info. of itself
pll upinfo [N]; // store info. of climbing up
int n, pa [N], layer [N], sz [N], done [N];
ll dis [__lg(N) + 1][N];
void init (int n)
   void init(int _n) {
      n = _n, layer[0] = -1;
      \begin{array}{ll} fill_n (pa + 1, n, 0), fill_n (done + 1, n, 0); \\ for (int i = 1; i <= n; ++i) G[i].clear(); \end{array}
   void add_edge(int a, int b, int w) {
      G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
   void get_cent(
      int\ u,\ int\ f\ ,\ int\ \&mx,\ int\ \&c\ ,\ int\ num)\ \{
      int mxsz = 0;
      sz[u] = 1;
      for (pll e : G[u])
          if (!done[e.X] && e.X != f) {
    get_cent(e.X, u, mx, c, num);
             sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
```

```
}
if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
  // if required, add self info or climbing info
  dis[layer[org]][u] = d;
  for ( dl e for )
    for (pll e : G[u])
if (!done[e.X] && e.X != f)
         dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    if (sz[e.X] > sz[c])
         \begin{array}{l} lc = cut(e.X, c, num - sz[c]);\\ else \ lc = cut(e.X, c, sz[e.X]); \end{array}
         upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
          a = pa[a], --ly)
       info[a].X += dis[ly][u], ++info[a].Y;
       if (pa[a])
         upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
  îl query(int u) {
    ll rt = 0;
    for (int a = u, ly = layer[a]; a;
          a = pa[a], --ly
       rt += info [a].X + info [a].Y * dis[ly][u];
       if (pa[a])
           upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
};
```

3.7 LiChaoST* [4a61ec]

```
ll m, k, id;
L(): id(-1) {}
 class LiChao { // maintain max
private:
 int n; vector<L> nodes;
  void insert(int l, int r, int rt, L ln) {
  int m = (l + r) >> 1;
    if (nodes[rt].id == -1)
      return nodes[rt] = ln, void();
    bool at Left = nodes[rt].at(1) < ln.at(1);
    if (nodes[rt].at(m) < ln.at(m))
      atLeft = 1, swap(nodes[rt], ln);
    if (r - l = 1) return;
if (atLeft) insert(l, m, rt << 1, ln);
    else insert (m, r, rt \ll 1 | 1, ln);
  11 query(int 1, int r, int rt, 11 x)
    int m = (1 + r) \gg 1; ll ret = -INF
    if (nodes[rt].id != -1) ret = nodes[rt].at(x);
    if (r - l = 1) return ret;
    if (x
        < m) return max(ret, query(1, m, rt << 1, x));
   return max(ret, query(m, r, rt \ll 1 | 1, x));
public:
 LiChao(int n_) : n(n_), nodes(n * 4)  {}
  void insert(\overline{L} ln) { insert(0, n, 1, ln); }
  ll query(ll x) { return query(0, n, 1, x); }
```

3.8 Link cut tree* [5f036a]

```
struct Splay { // xor-sum
  static Splay nil;
Splay *ch[2], *f;
```

```
int val, sum, rev, size;
  Splay (int
         _{\text{val}} = 0) : val(_{\text{val}}), sum(_{\text{val}}), rev(0), size(1)
    f = ch[0] = ch[1] = &nil;
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir()
  { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
     ch[d] = c;
     if (c != \&nil) c > f = this;
     pull();
  void give_tag(int r) {
     if (r) swap(ch[0], ch[1]), rev = 1;
  void push() {
  if (ch[0] != &nil) ch[0]->give_tag(rev);
  if (ch[1] != &nil) ch[1]->give_tag(rev);
     rev = 0:
  void pull() {
     // take care of the nil!
     Figure 1. Size + ch[1] - size + 1; sum = ch[0] - size + ch[1] - size + 1; sum = ch[0] - size + ch[1] - size + 1; if (ch[0] != &nil) ch[0] - f = this; if (ch[1] != &nil) ch[1] - f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate (Splay
                        *x) {
  Splay *p = x - > f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p\text{-}\!>\!setCh\left(x\text{-}\!>\!ch\left[\,!\,\dot{d}\,\right]\,,\;\;d\right);
  x->setCh(p, !d);
  p->pull(), x->pull();
void splay(Splay *x) {
  vector<Splay*> splayVec;
  for (Splay *q = x;; q = q->f) {
     splayVec.pb(q);
     if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
     if (x->f->isr()) rotate(x);
     else if (x->dir() = x->f->dir())
     rotate(x->f), rotate(x);
else rotate(x), rotate(x);
Splay* access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
     \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
  return q;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
  root_path(x), x->give_tag(1);
  x->push(), x->pull();
void split (Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link (Splay *x, Splay *y) {
  root_path(x), chroot(y);
x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y-size != 5) return;
  y->push();
  y->ch[0] = y->ch[0]->f = nil;
Splay* get_root(Splay *x)
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
    x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
  return get\_root(x) == get\_root(y);
```

```
Śplay* lca(Splay *x, Splay *y) {
  access(x), root_path(y);
  if (y->f = nil) return y;
  return y->f;
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
int query (Splay *x, Splay *y) {
  split(x, y);
  return y->sum;
3.9
     KDTree [74016d]
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl [maxn], yr [maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
  if (1 \rightleftharpoons r) return -1;
  function < bool (const point &, const point &)> f =
    [dep](const point &a, const point &b) {
  if (dep & 1) return a.x < b.x;</pre>
      else return a.y < b.y;
    };
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(1, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);

yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{rc}[m]]);
    yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds | | q.y > yr[o] + ds)
    return false:
  return true:
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
  const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
  if ((\text{dep \& 1}) \& \text{dep q.x} < p[o].x \mid )
    } else {
    if (\sim lc [o]) dfs (q, d, lc [o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
long long nearest (const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
  // namespace kdt
```

4 Flow/Matching

4.1 Bipartite Matching* [f07280]

```
National Taiwan University 8BQube
\begin{array}{lll} \textbf{struct} & \textbf{Bipartite\_Matching} ~\{~ / / ~\textbf{0-base} \\ & \textbf{int} & \textbf{mp}[N] ~, ~ \textbf{mg}[N] ~, ~ \textbf{dis} [N + 1] ~, ~ \textbf{cur} [N] ~, ~ 1 ~, ~ r ~; \end{array}
  vector < int > G[N + 1];
  bool dfs(int u) {
     for (int &i = cur[u]; i < SZ(G[u]); ++i) {
        int e = G[u][i];
        if (mq[e] == 1
               || (dis[mq[e]] = dis[u] + 1 \& dfs(mq[e]))|
          return mp[mq[e] = u] = e, 1;
     return dis[u] = -1, 0;
  bool bfs() {
     queue<int> q;
     fill_n(dis, l + 1, -1);
for (int i = 0; i < l; ++i)
if (!~mp[i])
          q.push(i), dis[i] = 0;
     while (!q.empty()) {
       int u = q. front();
        q.pop();
        for (int e : G[u])
          if (!~dis[mq[e]])
             q.\,push\,(mq[\,e\,]\,)\ ,\ dis\,[mq[\,e\,]\,]\ =\ dis\,[\,u\,]\ +\ 1;
     return dis[1] != -1;
  int matching() {
     int res = 0;
     fill_n(mp, l, -1), fill_n(mq, r, l);
     while (bfs()) {
       fill_n (cur, 1, 0);
        for (int i = 0; i < l; ++i)
          res += (! \sim mp[i] \&\& dfs(i));
     return res; // (i, mp[i] != -1)
  void add_edge(int s, int t) { G[s].pb(t); }
  void init(int _l, int _r) {
     l = _l, r = _r;
for (int i = 0; i <= 1; ++i)
       G[i].clear();
};
4.2 Kuhn Munkres* [edf909]
struct KM {
               // 0-base, maximum matching
  ll w[N][N], hl[N], hr[N], slk[N];
int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
bool vl[N], vr[N];
  void init(int _n) {
     n = _n;
for (int i = 0; i < n; ++i)
```

```
fill_n (w[i], n, -INF);
void add_edge(int a, int b, ll wei) {
  w[a][b] = wei;
fool Check(int x) {
  if (vl[x] = 1, ~fl[x])
    return vr[qu[qr++] = fl[x]] = 1;
  while (~x) swap(x, fr[fl[x] = pre[x]]);
  return 0;
void bfs(int s) {
  fill\_n\,(\,slk
       , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
  ql = qr = 0, qu[qr++] = s, vr[s] = 1;
  for (ll d;;) {
     while (q\hat{l} < qr)
for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] && slk
               [x] > = (d = hl[x] + hr[y] - w[x][y])) 
             if (pre[x] = y, d) slk[x] = d;
            else if (!Check(x)) return;
     d = INF;
     for (int x = 0; x < n; ++x)
       if (!vl[x] \& d > slk[x]) d = slk[x];
     for (int x = 0; x < n; ++x) {
  if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (\operatorname{vr}[x]) \operatorname{hr}[x] = d;
     for (int x = 0; x < n; ++x)
        if (!vl[x] && !slk[x] && !Check(x)) return;
```

4.3 MincostMaxflow* [47d9d2]

```
struct MinCostMaxFlow { // 0-base
   struct Edge {
      ll from, to, cap, flow, cost, rev;
   } *past[N];
   vector<Edge> G[N];
   int inq[N], n, s, t;
11 dis[N], up[N], pot[N];
bool BellmanFord() {
      fill_n(dis, n, \tilde{I}N\tilde{F}), fill_n(inq, n, 0);
      queue<int> q;
     auto relax = [&](int u, ll d, ll cap, Edge *e) {
  if (cap > 0 && dis[u] > d) {
           dis[u] = d, up[u] = cap, past[u] = e;
           if (!inq[u]) inq[u] = 1, q.push(u);
        }
      };
      relax(s, 0, INF, 0);
while (!q.empty()) {
        int u = q.front();
        q.pop(), inq[u] = 0;
for (auto &e : G[u]) {
           11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
                 (e.to, d2, min(up[u], e.cap - e.flow), &e);
        }
      return dis[t] != INF;
   void solve (int
     , int _t, ll &flow, ll &cost, bool neg = true) { s = \_s, t = \_t, flow = 0, cost = 0;
      if (neg) BellmanFord(), copy_n(dis, n, pot);
      for (; BellmanFord(); copy_n(dis, n, pot)) {
        for (int
        i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
        for (int i = t; past[i]; i = past[i]->from) {
          auto &e = *past[i];
           e.flow += up[t], G[e.to][e.rev].flow -= up[t];
        }
     }
   void init(int _n) {
     n = n, fill_n(pot, n, 0);
for (int i = 0; i < n; ++i) G[i].clear();
   void add_edge(ll a, ll b, ll cap, ll cost)
     \begin{array}{l} G[a].pb(Edge\{a,\ b,\ cap,\ 0,\ cost\ ,\ SZ(G[b])\});\\ G[b].pb(Edge\{b,\ a,\ 0,\ 0,\ -cost\ ,\ SZ(G[a])\ -\ 1\}); \end{array}
};
```

4.4 Maximum Simple Graph Matching*

```
struct Matching { // 0-base
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
        if (vis[x] == tk) return x;
        vis[x] = tk;
        x = Find(pre[match[x]]);
    }
  }
  void Blossom(int x, int y, int 1) {
    for (; Find(x) != 1; x = pre[y]) {
```

```
\begin{array}{l} pre\,[\,x\,] \,=\, y, \;\; y \,=\, match\,[\,x\,]\,; \\ if \;\; (\,s\,[\,y\,] \,\,=\,\, 1) \;\; q.\, push\,(\,y\,)\,, \;\; s\,[\,y\,] \,\,=\,\, 0\,; \\ for \;\; (\,int \;\; z\colon \,\{x, \;\,y\}) \;\; if \;\; (\,fa\,[\,z\,] \,\,=\,\, z\,) \;\; fa\,[\,z\,] \,\,=\,\, l\,; \end{array}
for (int x = q.front(); int u : G[x])
           if (s[u] = -1) {
if (pre[u] = x, s[u] = 1, match[u] = n) {
                   for (int a = u, b = x, last;
                          b != n; a = last, b = pre[a]
                             match[b], match[b] = a, match[a] = b;
                   return true;
           q.push(match[u]); s[match[u]] = 0;
} else if (!s[u] && Find(u) != Find(x)) {
                \begin{array}{ll} \text{int } l = LCA(u, x); \\ Blossom(x, u, l); \ Blossom(u, x, l); \end{array} 
   return false;
Matching(\, \underline{int} \, \, \underline{\hspace{1em}} n) \, : \, \, n(\underline{\hspace{1em}} n) \, , \, \, fa\,(n \, + \, 1) \, , \, \, s\,(n \, + \, 1) \, , \, \, vis \,
\begin{array}{c} (n+1)\,,\; pre\,(n+1,\;n)\,,\; match\,(n+1,\;n)\,,\; G(n) \;\; \{\}\\ void \;\; add\_edge\,(int\;\;u,\;\; int\;\;v) \end{array}
\{G[u].pb(v), G[v].pb(u); \}
int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
        if (match[x] = n) ans += Bfs(x);
    return ans:
} // match[x] == n means not matched
```

4.5 Maximum Weight Matching* [c80005]

```
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based</pre>
   struct edge { int u, v, w; }; int n, nx;
   vector<int> lab; vector<vector<edge>>> g;
   vector<int> slk, match, st, pa, S, vis;
   \begin{array}{l} \text{vector} < \text{int} > \text{ flo} \text{ , flo_from}; \text{ queue} < \text{int} > \text{ q}; \\ \text{WeightGraph}(\text{int n_}) : \text{n(n_}), \text{nx(n * 2), lab(nx + 1),} \\ \end{array}
      g(nx + 1, vector < edge > (nx + 1)), slk(nx + 1)
      flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
      match = st = pa = S = vis = slk;
     \label{eq:rep} \text{REP}(u, \ 1, \ n) \ \ \bar{\text{REP}}(v, \ 1, \ n) \ \ g[u][v] = \{u, \ v, \ 0\};
  { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; } void update_slk(int u, int x, int &s) { if (!s || E(g[u][x]) < E(g[s][x])) s = u; }
   void set_slk(int x) {
      slk[x] = 0;
     REP(u, 1, n)
         update_slk(u, x, slk[x]);
   void q_push(int x)
      if (x \le n) q.push(x);
      else for (int y : flo[x]) q_push(y);
   void set_st(int x, int b) {
      st[x] = b;
      if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
  auto it = find(ALL(f), xr);
  if (auto pr = it - f.begin(); pr % 2 == 1)
        reverse(1 + ALL(f)), it = f.end() - pr;
      auto res = vector(f.begin(), it);
     return f.erase(f.begin(), it), res;
   void set_match(int u, int v) {
     \operatorname{match}[\mathbf{u}] = g[\mathbf{u}][\mathbf{v}].\mathbf{v};
      if (u \le n) return;
     int xr = flo_from[u][g[u][v].u];

auto &f = flo[u], z = split_flo(f, xr);

REP(i, 0, SZ(z) - 1) set_match(z[i], z[i^1]);
     set_match(xr, v); f.insert(f.end(), ALL(z));
   void augment(int u, int v) {
      for (;;) {
```

```
int xnv = st[match[u]]; set_match(u, v);
if (!xnv) return;
     set_match(v = xnv, u = st[pa[xnv]]);
int lca(int u, int v) {
  static int t = 0; ++t;
   for (++t; u | | v; swap(u, v)) if (u) {
     if (vis[u] == t) return u;
     vis[u] = t, u = st[match[u]];
     if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int o, int v) {
  int b = find(n + 1 + ALL(st), 0) - begin(st);
  lab[b] = 0, S[b] = 0, match[b] = match[o];
   vector < int > f = \{o\};
   \  \, \text{for} \  \, (\, \text{int} \ t \ : \ \{u\,,\ v\}) \,\, \{ \,
     reverse(1 + ALL(f));
     for (int x = t, y; x != o; x = st[pa[y]])
        f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set_st(b, b);

REP(x, 1, nx) g[b][x].w = g[x][b].w = 0;

fill(ALL(flo_from[b]), 0);

for (int xs : flo[b]) {
    REP(x, 1, nx)
         \begin{array}{lll} & \text{if } (g[b][x].w = 0 \mid | E(g[xs][x]) < E(g[b][x])) \\ g[b][x] = g[xs][x], \ g[x][b] = g[x][xs]; \\ \end{array} 
    \begin{aligned} & \text{REP}(x, 1, n) \\ & \text{if } (\text{flo\_from}[xs][x]) \text{ flo\_from}[b][x] = xs; \end{aligned}
  set_slk(b);
void expand_blossom(int b) {
  for (int x : flo[b]) set_st(x, x);
int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  if (xs = -1) { xs = x; continue; } pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
     slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
   for (int x : flo[b])
     if(x = xr) S[x] = 1, pa[x] = pa[b];
     else S[x] = -1, set\_slk(x);
   st[b] = 0;
bool on_found_edge(const edge &e) {
  if (int u = st[e.u], v = st[e.v]; S[v] = -1) {
  int u = st[match[v]]; pa[v] = e.u; S[v] = 1;
  slk[v] = slk[nu] = S[nu] = 0; q_push(nu);
   else if (S[v] = 0)
     if (int \ o = lca(u, v)) add_blossom(u, o, v);
     else return augment(u, v), augment(v, u), true;
  return false;
bool matching() {
  fill(ALL(S), -1), fill(ALL(slk), 0);
   q = queue < int > ();

REP(x, 1, nx) \text{ if } (st[x] = x \&\& !match[x]) \\
pa[x] = S[x] = 0, q_push(x);

   if (q.empty()) return false;
  for (;;)
     while (SZ(q)) {
        int u = q.front(); q.pop();
        if (S[st[u]] = 1) continue;
        REP(v, 1, n)
           if (g[u][v].w > 0 & st[u] != st[v]) {
             if (E(g[u][v]) != 0)
                update_slk(u, st[v], slk[st[v]]);
                    (on_found_edge(g[u][v])) return true;
          }
      int d = INF;
    REP(b, n + 1, nx) if (st[b] = b && S[b] = 1)
d = min(d, lab[b] / 2);
     REP(x,\ 1,\ nx)
        if (int
              s = slk[x]; st[x] == x &  sk s &  S[x] <= 0
           d = min(d, E(g[s][x]) / (S[x] + 2));
    else if (S[st[u]] = 0) {
```

```
National Taiwan University 8BQube
             if (lab[u] <= d) return false;
            lab[u] -= d;
        \begin{array}{lll} \text{REP}(b, \ n+1, \ nx) & \text{if} & (\text{st}[b] \Longrightarrow b \ \&\& \ S[b] >= 0) \\ \text{lab}[b] & += d \ * (2 \ -4 \ * \ S[b]) \ ; \end{array} 
       REP(x, 1, nx)

if (int s = slk[x]; st[x] == x && \mathbb{R}
               s \&\& st[s] != x \&\& E(g[s][x]) == 0
             if (on_found_edge(g[s][x])) return true;
       REP(b, n + 1, nx)

if (st[b] = b && S[b] = 1 && lab[b] = 0)
            expand_blossom(b);
     return false;
  pair<ll, int> solve() {
     fill\left(ALL(match)\;,\;\;0\right);
    REP(u, 0, n) st [u] = u, flo [u]. clear();
     int w_max = 0;
    \label{eq:w_max} w\_max \, = \, \max(w\_max, \ g\left[\,u\,\right]\left[\,v\,\right].\, w) \ ;
     fill(ALL(lab), w_max);
     int n_matches = 0; ll tot_weight = 0;
     while (matching()) ++n_matches;
    REP(u, 1, n) if (match[u] \&\& match[u] < u)
       tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
  void add_edge(int u, int v, int w)
  \{\ g\,[\,u\,]\,[\,v\,]\,.\,w\,=\,g\,[\,v\,]\,[\,u\,]\,.\,w\,=\,w\,;\ \}
       SW-mincut [90bfe6]
4.6
```

```
\begin{array}{ll} int \ vst \, [MXN] \; , \ edge \, [MXN] \, [MXN] \; , \ wei \, [MXN] \; ; \\ void \ init \, (int \ n) \; \; \big\{ \end{array}
     REP fill_n (edge[i], n, 0);
   void addEdge(int u, int v, int w){
     edge[u][v] += w; edge[v][u] += w;
   int search(int &s, int &t, int n){
     fill_n(vst, n, 0), fill_n(wei, n, 0);
     s = t = -1;
     int mx, cur;
     for (int j = 0; j < n; ++j) {
        mx = -1, cur = 0;
        REP if (wei[i] > mx) cur = i, mx = wei[i];
        vst[cur] = 1, wei[cur] = -1;
       s = t; t = cur;
REP if (!vst[i]) wei[i] += edge[cur][i];
     return mx;
   int solve(int n) {
     int res = INF;
     for (int x, y; n > 1; n--){
         \begin{array}{l} res = min(res, search(x, y, n)); \\ REP \ edge[i][x] = (edge[x][i] += edge[y][i]); \\ \end{array} 
        REP {
           edge[y][i] = edge[n - 1][i];
edge[i][y] = edge[i][n - 1];
         // edge[y][y] = 0; 
     return res;
  }
} sw;
```

BoundedFlow*(Dinic*) [4ae8ab] 4.7

```
{\tt struct} \  \, {\tt BoundedFlow} \  \, \{\  \, //\  \, {\tt 0-base}
  struct edge {
    int to, cap, flow, rev;
  vector<edge> G[N];
  int n, s, t, dis[N], cur[N], cnt[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n + 2; ++i)
       G[i]. clear(), cnt[i] = 0;
  void add_edge(int u, int v, int lcap, int rcap) {
```

```
\begin{array}{l} \operatorname{cnt}[u] \mathrel{-=} \operatorname{lcap}, \; \operatorname{cnt}[v] \mathrel{+=} \operatorname{lcap}; \\ G[u].\operatorname{pb}(\operatorname{edge}\{v,\;\operatorname{rcap},\;\operatorname{lcap},\;\operatorname{SZ}(G[v])\}); \end{array}
      G[v].pb(edge\{u, 0, 0, SZ(G[u]) - 1\});
   void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, SZ(G[v])});
  G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
   int dfs(int u, int cap) {
  if (u = t || !cap) return cap;
       for (int &i = cur[u]; i < SZ(G[u]); ++i) {
         \stackrel{\cdot}{\mathrm{edge}} \& e = G[\underline{u}][i];
          if (dis[e.to] = dis[u] + 1 & e.cap! = e.flow) {
            int df = dfs(e.to, min(e.cap - e.flow, cap));
               e.flow += df, G[e.to][e.rev].flow -= df;
               return df;
            }
         }
      dis[u] = -1;
      {\tt return} \ 0;
   bool bfs() {
      fill_n(dis, n + 3, -1);
      queue<int> q;
      q.push(s), dis[s] = 0;
       while (!q.empty()) {
         int u = q.front();
         q.pop();
         for (edge &e : G[u])
            if (!~dis[e.to] && e.flow != e.cap)
               q.push(e.to), dis[e.to] = dis[u] + 1;
      return dis[t] != -1;
   int maxflow(int _s, int _t) {
      s = _s, t = _t;
      int flow = 0, df;
       while (bfs()) {
         \begin{array}{l} fill_n(cur, n + 3, 0); \\ while ((df = dfs(s, INF))) flow += df; \end{array}
      return flow;
   bool solve() {
      int sum = 0;
       for (int i = 0; i < n; ++i)
         if (cnt[i] > 0)
         \begin{array}{lll} & \text{add\_edge}(n+1,\ i,\ cnt[\,i\,])\ ,\ sum\ +=\ cnt[\,i\,]\,;\\ & \text{else if } (cnt[\,i\,]\ <\ 0)\ add\_edge(\,i\,,\ n+2\,,\ -cnt[\,i\,])\,; \end{array}
       if (sum != maxflow(n + 1, n + 2)) sum = -1;
       for (int i = 0; i < n; ++i)
         G[n + 1].pop_back(), G[i].pop_back();
else if (cnt[i] < 0)</pre>
            G[i].pop\_back(), G[n + 2].pop\_back();
      return sum != -1;
   int solve(int _s, int
      add_edge(_t, _s, INF);
if (!solve()) return -1; // invalid flow
       int x = G[\_t] . back() . flow;
      };
4.8 Gomory Hu tree* [5f2460]
```

```
MaxFlow Dinic;
\operatorname{int} \ g[MAXN];
void GomoryHu(int n) { // 0-base
   \begin{array}{ll} fill\_n\,(g\,,\ n\,,\ 0)\,;\\ for\ (int\ i\,=\,1;\ i\,<\,n;\ +\!\!+\!\!i\,)\ \{ \end{array}
       Dinic.reset();
       {\rm add\_edge}(i\,,\,g[\,i\,]\,,\,\,{\rm Dinic.maxflow}(i\,,\,\,g[\,i\,])\,)\,;
       for (int j = i + 1; j \le n; ++j)
          if (g[j] = g[i] \& cond \sim Dinic.dis[j])
             g[j] = i;
}
```

Minimum Cost Circulation* [cb40c6]

```
struct MinCostCirculation { // 0-base
  struct Edge {
    ll from, to, cap, fcap, flow, cost, rev;
```

```
} *past[N]:
  vector<Edge> G[N];
  11 dis[N], inq[N], n;
void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
auto relax = [&](int u, ll d, Edge *e) {
       if (dis[u] > d) {
dis[u] = d, past[u] = e;}
         if (!inq[u]) inq[u] = 1, q.push(u);
       }
    relax(s, 0, 0);
    while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
            relax(e.to, dis[u] + e.cost, \&e);
    }
  }
  void try_edge(Edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
         \begin{tabular}{ll} $i=cur.from; $past[i]$; $i=past[i]$->from) { & auto &e = *past[i]$; $++e.flow, $--G[e.to][e.rev].flow$; } \end{tabular}
       }
    ++cur.cap;
  void solve(int mxlg) {
    for (int b = mxlg; b >= 0; --b) {
  for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
       e.cap *= 2, e.flow *= 2;
for (int i = 0; i < n; ++i)
         for (auto &e : G[i])
            if (e.fcap >> b & 1)
              try_edge(e);
  }
  void init (int \underline{n}) { n = \underline{n};
    for (int i = 0; i < n; ++i) G[i].clear();
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge
    } mcmf; // O(VE * ElogC)
```

4.10 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
- 1. Construct supersource S and sink T.
- 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - − To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex coverfrom maximum matching M on bipartite $\operatorname{graph}(X,Y)$
 - 1. Redirect every edge: $y \to x$ if $(x,y) \in M, x \to y$ otherwise.
 - 2. DFS from unmatched vertices in X.
 - 3. $x \in X$ is chosen iff x is unvisited.
 - 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
- 1. Consruct super source S and $\sinh T$
- 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
- 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v) > 0, connect $S \rightarrow v$ with (cost, cap) = (0, d(v))

- 5. For each vertex v with d(v)<0, connect $v\to T$ with (cost,cap)=(0,-d(v))
- 6. Flowfrom S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer T
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - $3. \ \ The {\it mincut} is equivalent to the {\it maximum} profit of a subset of projects.$
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_{uv}
 - 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \ge 0$$

5 String

5.1 KMP [9e1cd1]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
  }
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  }
  return ans;
}</pre>
```

5.2 Z-value* [e2dc6f]

5.3 Manacher* [bfe74e]

```
5.4 SAIS* [e9a275]
auto sais (const auto &s) {
  const int n = SZ(s), z = ranges::max(s) + 1;
  if (n = 1) return vector\{0\};
vector\{\text{int} > c(z); for (\text{int } x : s) ++c[x];
  partial\_sum(ALL(c), begin(c));
  t[i] = (
          s[i] = s[i+1] ? t[i+1] : s[i] < s[i+1]);
  auto is_lms = views::filter([&t](int x) {
     return x && t[x] && !t[x - 1];
  });
  auto induce = [\&] {
     for (auto x = c; int y : sa)
     if (y-) if (!t[y]) sa[x[s[y]-1]++]=y;
for (auto\ x=c;\ int\ y:\ sa\ |\ views::reverse)
        if (y--) if (t[y]) sa[--x[s[y]]] = y;
  vector < int > lms, q(n); lms.reserve(n);
  for (auto x = c; int i : I | is_lms)
q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
  induce(); vector < int > ns(SZ(lms));
  for (int j = -1, nz = 0; int i : sa \mid is\_lms) {
     if (j >= 0) {
        int len = min({n - i, n - j, lms[q[i] + 1] - i});
       ns[q[i]] = nx += lexicographical_compare(
begin(s) + j, begin(s) + j + len,
begin(s) + i, begin(s) + i + len);
     j = i;
  fill(ALL(sa), 0); auto nsa = sais(ns);
for (auto x = c; int y : nsa | views::reverse)
     y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
// sa[i]: sa[i]-th suffix
       is the i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
  int n; vector < int > sa, hi, ra;
  Suffix
     (const auto &_s, int _n) : n(_n), hi(n), ra(n) {
vector<int> s(n + 1); // s[n] = 0;
copy_n(_s, n, begin(s)); // _s shouldn't contain 0
sa = sais(s); sa.erase(sa.begin());
for (int i = 0; i < n; ++i) ra[sa[i]] = i;
for (int i = 0, b = 0; i < n; ++i) f
     for (int i = 0, h = 0; i < n; ++i)
        if (!ra[i]) { h = 0; continue; }
        for (int j = sa[ra[i] - 1]; max
             (i, j) + h < n &  s[i + h] = s[j + h];) + h;
       hi[ra[i]] = h ? h-- : 0;
  }
};
```

5.5 Aho-Corasick Automatan* [91c6c0]

```
struct AC_Automatan {
  int nx[len][sigma], fl[len], cnt[len], ord[len], top;
  int rnx[len][sigma]; // node actually be reached
  int newnode() {
    fill\_n\left(nx[top]\;,\;sigma\,,\;-1\right);
    return top++;
  void init() \{ top = 1, newnode(); \}
  int input(string &s) {
    int X = 1;
    for (char c : s) {    if (!\simnx[X][c - 'A']) nx[X][c - 'A'] = newnode();    X = nx[X][c - 'A'];
    return X; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(1), fl[1] = 0;
for (int t = 0; !q.empty(); ) {
      int R = q.front();
       q.pop(), ord[t++] = R;
       for (int i =
                      0; i < sigma; ++i)
         if (~nx[R][i])
           int X = rnx[R][i] = nx[R][i], Z = fl[R];
           for (; Z &  : \neg nx[Z][i]; ) Z = fl[Z];
```

```
fl[X] = Z ? nx[Z][i] : 1, q.push(X);
}
else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
}

void solve() {
  for (int i = top - 2; i > 0; --i)
     cnt[fl[ord[i]]] += cnt[ord[i]];
}
} ac;
```

5.6 Smallest Rotation [e74dc0]

```
\begin{array}{l} string \ mcp(string \ s) \ \{ \\ int \ n = SZ(s) \,, \ i = 0 \,, \ j = 1; \\ s += s \,; \\ while \ (i < n \ \&\& \ j < n) \ \{ \\ int \ k = 0 \,; \\ while \ (k < n \ \&\& \ s[i + k] \Longrightarrow s[j + k]) \ ++k; \\ if \ (s[i + k] <= s[j + k]) \ j += k + 1; \\ else \ i += k + 1 \,; \\ if \ (i \Longrightarrow j) \ ++j \,; \\ \} \\ int \ ans = i < n \ ? \ i \ : j \,; \\ return \ s.substr(ans, n) \,; \\ \} \end{array}
```

5.7 De Bruijn sequence* [f601c2]

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  if (ptr >= L) return;
    if (t > N) {
       if (N % p) return;
      for (int i = 1; i \le p \&\& ptr < L; ++i)
        \operatorname{out}[\operatorname{ptr}++] = \operatorname{buf}[i];
    } else
      buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
       for (int j = buf[t - p] + 1; j < C; ++j)
        buf[t] = j, dfs(out, t + 1, t, ptr);
  void solve(int _c, int _n, int _k, int *out) {
    int p = 0;
    C = c, N = n, K = k, L = N + K - 1; dfs(out, 1, 1, p);
    if (p < L) fill (out + p, out + L, 0);
} dbs;
```

5.8 Extended SAM* [58fa19]

```
struct exSAM {
  int len[N *
  int len [N * 2], link [N * 2]; // maxlength, suflink int next [N * 2] [CNUM], tot; // [0, tot), root = 0 int lenSorted [N * 2]; // topo. order int cnt [N * 2]; // occurence int newpools [(
   int newnode() {
      fill_n (next[tot], CNUM, 0);
      len[tot] = cnt[tot] = link[tot] = 0;
      return tot++;
   void init() { tot = 0, newnode(), link[0] = -1; }
   int insertSAM(int last, int c) {
  int cur = next[last][c];
      len[cur] = len[last] + 1;
      int p = link[last];
      while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
if (p = -1) return link[cur] = 0, cur;
      int q = next[p][c];
      if (len
            [p] + 1 = len[q] return link[cur] = q, cur;
      int clone = newnode();
      \quad \text{for (int } i = 0; i < \overleftarrow{CNUM}; ++i)
               clone][i] = len[next[q][i]] ? next[q][i] : 0;
      len[clone] = len[p] + 1;
      while (p != -1 && next[p][c] == q)
        next \left[ p \right] \left[ \, c \, \right] \, = \, clone \, , \  \, p \, = \, link \left[ \, p \, \right] ; \label{eq:clone_p}
      link[link[cur] = clone] = link[q];
      link[q] = clone;
      return cur;
   void insert(const string &s) {
```

```
const string u = s.substr(0, nu), v = s.substr(nu),
     int cur = 0:
     for (auto ch : s) {
                                                                                     ru(u.rbegin
       int & nxt = next[cur][int(ch - 'a')];
                                                                                           (), u.rend()), rv(v.rbegin(), v.rend());
                                                                              main_lorentz(u, sft), main_lorentz(v, sft + nu);
const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u)
        if (!nxt) nxt = newnode();
       \operatorname{cnt}\left[\operatorname{cur} = \operatorname{nxt}\right] += 1;
                                                                              z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v);
auto get_z = [](const vector<int> &z, int i) {
    }
  }
  void build() {
     queue<int> q;
                                                                                       (0 \le i \text{ and } i \le (int)z.size()) ? z[i] : 0; };
     q.push(0);
                                                                              auto add_rep
     while (!q.empty()) {
                                                                                   = [\&](bool left, int c, int l, int k1, int k2) {
       int cur = q.front();
                                                                                 const
        q.pop();
                                                                                       int L = max(1, 1 - k2), R = min(1 - left, k1);
       for (int i = 0; i < CNUM; ++i)
if (next[cur][i])
                                                                                 if (L > R) return;
                                                                                if (left)
                                                                                rep[l].emplace_back(sft + c - R, sft + c - L);
else rep[l].emplace_back
            q.push(insertSAM(cur, i));
     vector < int > lc(tot);
                                                                                      (sft + c - R - l + 1, sft + c - L - l + 1);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];
     partial_sum(ALL(lc), lc.begin());
                                                                              for (int cntr = 0; cntr < n; cntr++) {
                                                                                int 1, k1, k2;
     for (int i
          = \ 1; \ i < tot; \ +\!\!+\! i \,) \ lenSorted[\, -\! -\! lc \, [\, len \, [\, i \, ]\, ]\,] \ = \ i \,;
                                                                                 if (cntr < nu) {
                                                                                   l = nu - cntr;
  void solve() {
  for (int i = tot - 2; i >= 0; --i)
                                                                                   k1 = get_z(z1, nu - cntr);
                                                                                   k2 = get_z(z^2, nv + 1 + cntr);
       cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
                                                                                } else {
                                                                                   l = cntr - nu + 1;
};
                                                                                   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
                                                                                   k2 = get_z(z4, (cntr - nu) + 1);
     PalTree* [675736]
5.9
                                                                                 if (k1 + k2 >= 1)
{\tt struct palindromic\_tree} \ \{
                                                                                   add_rep(cntr < nu, cntr, l, k1, k2);
  struct node {
     int next[26], fail, len;
                                                                          | \} // p \in [1, r] \Rightarrow s[p, p+i) = s[p+i, p+2i)
     6 Math
       for (int i = 0; i < 26; ++i) next[i] = 0;
                                                                           6.1 ax+by=gcd(only exgcd *) [5fef50]
  };
                                                                           pll exgcd(ll a, ll b) {
  vector<node> St:
                                                                              if (b = 0) return pll(1, 0);
  vector<char> s;
                                                                              ll p = a / b;
  int last, n;
                                                                              pll q = exgcd(b, a \% b);
  palindromic\_tree() : St(2), last(1), n(0) {
                                                                              return pll(q.Y, q.X - q.Y * p);
    St[0].fail = 1, St[1].len = -1, s.pb(-1);
                                                                           /* ax+by=res, let x be minimum non-negative
  inline void clear() {
                                                                          \begin{array}{l} (g, \ p = \gcd(a, \ b), \ \exp(d(a, \ b)) \ * \ res \ / \ g \\ (g, \ p = \gcd(a, \ b), \ \exp(d(a, \ b)) \ * \ res \ / \ g \\ (g, \ p = \gcd(a, \ b), \ \exp(d(a, \ b)) \ * \ res \ / \ g \\ (g, \ p = (b, X \ / \ b) \ / \ g \ - 1) \ / \ (b \ / \ g) \\ (g, \ p = (b, A \ / \ g)) \ + (b, A \ g) \ * \ t \ */ \end{array}
     St.clear(), s.clear(), last = 1, n = 0; St.pb(0), St.pb(-1);
     St[0]. fail = 1, s.pb(-1);
  inline int get_fail(int x) {
                                                                           6.2 Floor and Ceil [1ffa73]
     while (s[n - St[x].len - 1] != s[n])
       x = \hat{S}t[x].fail;
                                                                           int floor(int a, int b)
     return x;
                                                                           { return a / b - (a % b && (a < 0) \hat{} (b < 0)); } int ceil(int a, int b)
  inline void add(int c) {
   s.push_back(c -= 'a'), ++n;
                                                                           \{ \text{ return } a \ / \ b + (a \% b \&\& (a < 0) ^ (b > 0)); \} 
     int cur = get_fail(last);
     if (!St[cur].next[c]) {
  int now = SZ(St);
                                                                           6.3 Floor Enumeration [67ad61]
        St.pb(St[cur].len + 2);
                                                                           // enumerating x = floor(n / i), [l, r]
                                                                           for (int l = 1, r; l <= n; l = r + 1) {
int x = n / l;
       St [now]. fail
          St [get_fail(St [cur].fail)].next[c];
        St[cur].next[c] = now;
                                                                              r = n / x;
       St[now].num = St[St[now].fail].num + 1;
     last = St[cur].next[c], ++St[last].cnt;
                                                                           6.4 Mod Min [038fef]
  inline void count() { // counting cnt
                                                                           // \min\{k \mid 1 \le ((ak) \mod m) \le r\}, no solution -> -1 ll \max_{\min}(ll \ a, \ ll \ m, \ ll \ l, \ ll \ r)  { if (a == 0) return l \ ? \ -1 : \ 0; if (ll \ k = (l + a \ - \ 1) \ / \ a; \ k \ * \ a <= \ r)
     auto i = St.rbegin();
     for (; i != St.rend(); ++i) {
       St[i->fail].cnt += i->cnt;
                                                                                return k;
                                                                              11 b = m / a, c = m \% a;
  inline int size() { // The number of diff. pal.
                                                                              if (11 y = mod_min(c, a, a - r % a, a - 1 % a))
return (1 + y * c + a - 1) / a + y * b;
```

5.10 Main Lorentz [eaf279]

return SZ(St) - 2;

};

```
vector<pair<int, int>>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
 const int n = s.size();
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
```

6.5 Linear Mod Inverse [aa1426]

```
inv[1] = 1;
for ( int i = 2; i
     <= N; ++i ) inv[i] = ((mod-mod/i)*inv[mod%i])%mod;
```

6.6 Linear Filter Mu [663a36]

6.7 Gaussian integer gcd [4fcbff]

6.8 GaussElimination [c016c9]

```
\begin{tabular}{ll} \begin{tabular}{ll} void $GAS(V<V<double>>&vc) & \{ \end{tabular}
      int len = vc.size();
      for (int i = 0; i < len; ++i)
            int idx = find_if(vc.begin()+i, vc.end(),[&](
                  if ( idx = len ) continue
            if( i != idx ) swap( vc[idx], vc[i] );
double pivot = vc[i][i];
            for_each( vc[i].begin(), vc
            for_each( vc[i].bcgin(), vc
  [i].end(), [&]( auto &a ) { a/=pivot; } );
for( int j = 0; j < len; ++j ) {
  if( i == j ) continue;
  if( vc[j][i]!= 0 ) {
     double mul = vc[j][i]/vc[i][i];
     trueform( va[i] bogin() va[i] end</pre>
                         transform(vc[j].begin(), vc[j].end
                               (), vc[i].begin(), vc[j].begin(),
                                     [&](auto &a, auto &b) {
                                     return a-b*mul;
                                     });
                  }
            }
     }
};
```

6.9 Miller Rabin* [14b81a]

6.10 Simultaneous Equations [21b2e1]

```
| struct matrix { //m variables, n equations int n, m; fraction M[MAXN] [MAXN + 1], sol [MAXN]; int solve() { //-1: inconsistent, >= 0: rank for (int i = 0; i < n; ++i) { int piv = 0; while (piv < m && !M[i][piv].n) ++piv; if (piv = m) continue; for (int j = 0; j < n; ++j) { if (i == j) continue;
```

6.11 Pollard Rho* [fff0fc]

6.12 Simplex Algorithm [40618e]

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a [MAXN] [MAXM] , b [MAXN] , c [MAXM];
double d [MAXN] [MAXM] , x [MAXM];
int ix [MAXN + MAXM]; //!!! array all indexed from 0
// max{cx} subject to {Ax=b,x>=0}
// n: constraints, m: vars !!!
   x[] is the optimal solution vector
// usage :
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
   fill_n(d[n], m + 1, 0);
   fill_n(d[n+1], m+1, 0);
  iota(ix, ix + n + m, 0);
  \  \  \, \text{int}\  \  \, r\,=\,n\,,\  \  \, s\,=\,m\,\,\,\text{-}\,\,\,1\,;
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];
     d[i][m-1] = 1;
     d[i][m] = b[i];
     if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);

d[n + 1][m - 1] = -1;
  for (double dd;; ) {
     if'(r < n) {
        swap(ix[s], ix[r+m]);

d[r][s] = 1.0 / d[r][s];
       }
     r = s = -1;
     for (int j = 0; j < m; ++j)

if (s < 0 || ix[s] > ix[j]) {

if (d[n + 1][j] > eps ||
                (d[n + 1][j] > -eps \& d[n][j] > eps))
     if'(s < 0) break;
     for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
        if(r < 0 | |
```

```
(dd = d[r][m]
                    /d[r][s] - d[i][m] / d[i][s]) < -eps | |
            (dd < eps & ix[r+m] > ix[i+m]))
   if (r < 0) return -1; // not bounded
if (d[n + 1][m] < -eps) return -1; // not executable
double ans = 0;
fill_n(x, m, 0);
for (int i = m; i <
   \begin{array}{ll} (in \ i - in, \ i - in) & \text{if } (ix \ i \ i \ i - in) & \text{if } (ix \ [i \ ] < m - 1) \\ ans & \text{if } (ix \ [i \ ] < m - 1) & \text{if } (ix \ [i \ ] ; \end{array}
      x[ix[i]] = d[i-m][m];
}
return ans;
```

6.12.1 Construction

Primal	Dual
Maximize $c^{\intercal}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \geq c, y \geq 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T} y$ s.t. $A^{T} y = c, y \ge 0$
Maximize $c^{T}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{T}y$ s.t. $A^{T}y \ge c$

 $\overline{\mathbf{x}}$ and $\overline{\mathbf{y}}$ are optimalified only if for all $i \in [1, n]$, either $\overline{x}_i = 0$ or $\sum_{j=1}^m A_{ji} \overline{y}_j = c_i$ holds and for all $i \in [1, m]$ either $\bar{y}_i = 0$ or $\sum_{j=1}^n A_{ij} \bar{x}_j = b_j$ holds.

- 1. In case of minimization, let $c_i' = -c_i$ 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \rightarrow \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$

- $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$ $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.13 chineseRemainder [fe9f25]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
 ll g = \gcd(m1, m2);
 if ((x2 - x1) % g) return -1; // no sol m1 /= g; m2 /= g;
  pll p = exgcd(m1, m2);
 ll lcm = m1 * m2 * g;
ll res = p.first * (x2 - x1) * m1 + x1;
  // be careful with overflow
  return (res % lcm + lcm) % lcm;
```

6.14 Factorial without prime factor* [dcffcb]

```
O(p^k + \log^2 n), pk = p^k
ll prod [MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
   \operatorname{prod}[0] = 1;
  if or (int i = 1; i <= pk; ++i)
   if (i % p) prod[i] = prod[i - 1] * i % pk;
   else prod[i] = prod[i - 1];</pre>
   11 \text{ rt} = 1;
   for (; n; n /= p) {
  rt = rt * mpow(prod[pk], n / pk, pk) % pk;
     rt = rt * prod[n \% pk] \% pk;
} // (n! without factor p) % p^k
```

Discrete Log* [ba4ac0]

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered\_map < int, int > p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {
     \begin{aligned} p[y] &= i; \\ y &= 1LL * y * x % m; \\ b &= 1LL * b * x % m; \end{aligned} 
  for (int i = 0; i < m + 10; i += kStep) {
s = 1LL * s * b % m;
     if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m = 1) return 0;
  int s = 1:
  for (int i = 0; i < 100; ++i) {
     if (s == y) return i;
s = 1LL * s * x % m;
```

```
if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
}
```

6.16 Berlekamp Massey [9380b8]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
   vector < T > d(SZ(output) + 1), me, he;
   for (int j = 0, i = 1; i <= SZ(output); ++i) {
  for (int j = 0; j < SZ(me); ++j)
   d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;
  if (me.empty()) {
    me.resize(f - i):
         me. resize(f = i);
         continue;
     vector < T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
for (T x : he) o.pb(x * k);
      o.resize(max(SZ(o), SZ(me)));
      for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
      if (i - f + SZ(he)) = SZ(me) he = me, f = i;
      me = o;
   return me;
```

6.17Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 98789101 987777733 999991921 1010101333
     1010102101 \ 1000000000039 \ 100000000000037
     2305843009213693951 \ \ 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

Theorem 6.18

• Cramer'srule

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i), L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the $\max \operatorname{maximum} \operatorname{matching} \operatorname{on} G.$

- Cayley's Formula
 - Given a degree sequence $d_1, d_2, ..., d_n$ for each labeled vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that $\operatorname{vertex} 1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}.$
- Erdős–Gallaitheorem

A sequence of nonnegative integers $d_1 \ge \cdots \ge d_n$ can be represented as the degreesequenceofafinitesimplegraphon n vertices if and only if $d_1 + \cdots + d_n$

is even and $\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i,k)$ holds for every $1 \le k \le n$.

 $Gale\!-\!Rysertheorem$

A pair of sequences of nonnegative integers $a_1 \ge \cdots \ge a_n$ and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^{n} a_i = \sum_{i=1}^{n} b_i$ and $\sum_{i=1}^{k} a_i \leq \sum_{i=1}^{n} \min(b_i, k)$ holds for every $1 \le k \le n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \geq \cdots \geq a_n$ is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1) + \sum_{i=1, i+1}^n \min(b_i, k) \operatorname{holds} \text{ for every } 1 \leq k \leq n.$

Pick's theorem

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1.$

• Möbius inversion formula

$$\begin{array}{ll} - & f(n) = \sum_{d \mid n} g(d) \Leftrightarrow g(n) = \sum_{d \mid n} \mu(d) f(\frac{n}{d}) \\ - & f(n) = \sum_{n \mid d} g(d) \Leftrightarrow g(n) = \sum_{n \mid d} \mu(\frac{d}{n}) f(d) \end{array}$$

- Sphericalcap
 - $\ A \, portion of a \, sphere \, cut \, off \, by \, a \, plane.$
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ :
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \sin \theta)$ $\cos\theta)^2/3$.
 - Area = $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2 (1 \cos\theta)$.
- Lagrange multiplier
 - Optimize $f(x_1,...,x_n)$ when k constraints $g_i(x_1,...,x_n) = 0$.
 - Lagrangian function $\mathcal{L}(x_1, \dots, x_n, \lambda_1, \dots, \lambda_k) = f(x_1, \dots, x_n)$ $\sum_{i=1}^{k} \lambda_i g_i(x_1, ..., x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines
 - Line 1: $v_1 = p_1 + t_1 d_1$
 - Line 2: $v_2 = \bar{p_2} + t_2 d_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$

 - $\mathbf{n}_1 = \mathbf{d}_1 \times \mathbf{n} \\ \mathbf{n}_2 = \mathbf{d}_2 \times \mathbf{n}$
- Derivatives/Integrals

Derivatives/Integrals Integrals Integration by parts:
$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$

$$\begin{vmatrix} \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx}\tan x = 1 + \tan^2 x \\ \int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \end{vmatrix} \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1)$$

$$\int \sqrt{a^2 + x^2} = \frac{1}{2}\left(x\sqrt{a^2 + x^2} + a^2 \sinh(x/a)\right)$$

• Spherical Coordinate

$$(x,y,z) = (r\sin\theta\cos\phi, r\sin\theta\sin\phi, r\cos\theta)$$

$$(r,\theta,\phi) = (\sqrt{x^2 + y^2 + z^2}, a\cos(z/\sqrt{x^2 + y^2 + z^2}), a\tan(y,x))$$

• Rotation Matrix

6.19Estimation

 $n \hspace{0.2cm} | \hspace{0.04cm} 2345 \hspace{0.1cm} 6 \hspace{0.1cm} 7 \hspace{0.1cm} 8 \hspace{0.1cm} 9 \hspace{0.1cm} 20 \hspace{0.1cm} 30 \hspace{0.1cm} 40 \hspace{0.1cm} 50 \hspace{0.1cm} 100$

p(n) 23571115223062756044e42e52e8

 $n \hspace{0.1cm} | 1001e31e6 \hspace{0.1cm} 1e9 \hspace{0.1cm} 1e12 \hspace{0.1cm} 1e15 \hspace{0.1cm} 1e18$

d(i) 12 32 2401344672026880103680

 $n \mid 123456789$ 10 11 12 13 14 15 $\frac{\binom{2n}{n}}{2} \, 2\, 6\, 20\, 70\, 252\, 924\, 3432\, 12870\, 48620\, 184756\, 7e5\, 2e6\, 1e7\, 4e7\, 1.5e8$ n 2 3 4 5 6 7 8 9 10 11 12 13 $B_n | 2515522038774140211471159757e54e63e7$

6.20 Euclidean Algorithms

- $m = |\frac{an+b}{c}|$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c,c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} a,b,c,n) &= \sum_{i=0} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \operatorname{mod} c,b \operatorname{mod} c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \begin{cases} \frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1)) \\ -h(c,c-b-1,a,m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ - 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

General Purpose Numbers

Bernoulli numbers
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k}$$
Strick a number of the article in Positivians of a distribution of a dis

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {i \choose i} i^{n}$$

$$x^{n} = \sum_{i=0}^{n} S(n,i)(x)_{i}$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
• Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. kj:ss.t. $\pi(j) > \pi(j+1), k+1j$:ss.t. $\pi(j) \ge j$, kj:ss.t. $\pi(j) > j$. E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)E(n,0) = E(n,n-1) = 1

$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$ 6.22 Tips for Generating Functions

- Ordinary Generating Function $A(x) = \sum_{i>0} a_i x^i$
 - $-A(rx) \Rightarrow r^n a_n$
 - $A(x)+B(x) \Rightarrow a_n+b_n$
 - $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
 - $-A(x)^{k} \Rightarrow \sum_{i_{1}+i_{2}+\cdots+i_{k}=n} a_{i_{1}} a_{i_{2}} \dots a_{i_{k}}$
 - $-xA(x)' \Rightarrow na_n$
 - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$
 - $-A(x)+B(x) \Rightarrow a_n+b_n$

 - $-A^{(k)}(x) \Rightarrow a_{n+k}$ $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$ $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n}^{n} {n \choose i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
 - $-xA(x) \Rightarrow na_n$
- Special Generating Function
 - $-(1+x)^n = \sum_{i\geq 0} {n \choose i} x^i$
 - $-\frac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{i}{n-1} x^i$ $\mathbf{Polynomial}$

Fast Fourier Transform [9fec80]

```
const int maxn = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
     for (int i = 0; i \le maxn)
           ; ++i) \text{ omega}[i] = \exp(i * 2 * pi / maxn * I);
void bin(vector<cplx> &a, int n) {
     int lg;
    \begin{array}{ll} \text{for (lg = 0; (1 << lg) < n; ++lg); --lg;} \\ \text{vector} < \text{cplx} > \text{tmp(n);} \end{array}
     for (int i = 0; i < n; ++i)
          int to = 0;
          for (int j = 0; (1 << j) <
n; ++j) to |= (((i >> j) & 1) << (lg - j));
          tmp[to] = a[i];
```

```
}
     for (int i = 0; i < n; ++i) a[i] = tmp[i];
                                                                        7.3 Fast Walsh Transform* [36c9f5]
}
                                                                       /* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
void fft(vector<cplx> &a, int n) {
                                                                        xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
     bin(a, n);
     for (int step = 2; step <= n; step <<= 1) {
                                                                       void fwt(int *a, int n, int op) { //or
for (int L = 2; L <= n; L <<= 1)
for (int i = 0; i < n; i += L)</pre>
          int to = step \gg 1;
          for (int i = 0; i < n; i += step) {
               for (int k = 0; k < to; ++k) {
                    cplx x = a[i
                                                                               for (int j = i; j < i + (L >> 1); ++j)

a[j + (L >> 1)] += a[j] * op;
                        + to + k] * omega[maxn / step * k];
                    a[i + to + k] = a[i + k] - x;
                    a[i + k] += x;
                                                                        const int N = 21;
               }
                                                                        int f
         }
                                                                            N[1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
     }
                                                                        void
}
                                                                            subset_convolution(int *a, int *b, int *c, int L) {
                                                                           // c_k = \sum_{i = k, i \& j = 0} a_i * b_j
void ifft (vector < cplx > &a, int n) {
                                                                          int n = 1 \ll L;
     fft(a, n);
                                                                          for (int i = 1; i < n; +++i)
     reverse(a.begin() + 1, a.end());
                                                                            ct[i] = ct[i \& (i - 1)] + 1;
                                                                           for (int i = 0; i < n; ++i)
     for (int i = 0; i < n; ++i) a[i] /= n;
                                                                           \begin{array}{l} f[\,ct\,[\,i\,]\,][\,i\,] = a\,[\,i\,]\,,\; g[\,ct\,[\,i\,]\,][\,i\,] = b\,[\,i\,]\,;\\ for\;\; (\,int\,\;i\,=\,0;\;\;i\,<=\,L\,;\;+\!\!+\!\!i\,) \end{array}
vector<int> multiply(const vector<
                                                                             fwt(f[i]\,,\,n,\,1)\,,\,fwt(g[i]\,,\,n,\,1)\,;
     int> &a, const vector<int> &b, bool trim = false) {
                                                                           for (int i = 0; i \ll L; ++i)
                                                                            for (int j = 0; j <= i; ++j)

for (int x = 0; x < n; ++x)

h[i][x] += f[j][x] * g[i - j][x];
     int d = 1;
          (d < max(a.size(), b.size())) d <<= 1; d <<= 1;
      vector <\! cplx\! > \, pa(d)\,, \ pb(d)\,;
                                                                          for (int i = 0; i <= L; ++i)
fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
     for (int i
           = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
     for (int i
                                                                            c[i] = h[ct[i]][i];
           = 0; i < b.size(); ++i) pb[i] = cplx(b[i], 0);
     fft\left(pa\,,\ d\right);\ fft\left(pb\,,\ d\right);
                                                                        7.4 Polynomial Operation [37b8c7]
     for (int i = 0; i < d; ++i) pa[i] *= pb[i];
     ifft (pa, d);
     vector < int > r(d);
                                                                       for (int
           i = 0; i < d; ++i) r[i] = round(pa[i].real());
     if (trim)
                                                                          using vector<ll>>::vector;
          while (r.size() \&\& r.back() == 0) r.pop\_back();
                                                                           static NTT<MAXN, P, RT> ntt;
     return r;
                                                                          int n() const { return (int)size(); } // n() >= 1
Poly(const Poly &p, int m): vector<ll>(m) {
}
                                                                             copy\_n(p.data(), min(p.n(), m), data());
 Prime
              Root
                      Prime
                                    Root
 7681
              17
                      167772161
                                    3
                                                                          Poly& irev()
 12289
                      104857601
              11
                                                                          { return reverse(data(), data() + n()), *this; } Poly& isz(int m) { return resize(m), *this; } Poly& iadd(const Poly &rhs) { // n() == rhs.n()
 40961
              3
                      985661441
                      998244353
 65537
              3
                                    3
 786433
              10
                      1107296257
                                    10
                                                                            fi(0, n()) if
   (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
 5767169
              3
                      2013265921
                                    31
 7340033
                      2810183681
                                    11
 23068673
              3
                      2885681153
                                    3
  469762049
                      605028353
                                                                          Poly& imul(ll k) {
       Number Theory Transform* [eleb36]
7.2
                                                                             fi(0, n()) (*this)[i] = (*this)[i] * k \% P;
vector<int> omega;
                                                                             return *this;
void Init() {
  omega. resize(kN + 1);
                                                                          Poly Mul(const Poly &rhs) const {
  long long x = fpow(kRoot, (Mod - 1) / kN);
                                                                             int m = 1:
  omega[0] = 1;
                                                                             while (m < n() + rhs.n() - 1) m <<= 1;
   for (int i = 1; i \le kN; ++i)
                                                                             Poly X(*this, m), Y(rhs, m);
     omega[i] = 1LL * omega[i - 1] * x % kMod;
                                                                             ntt\left(X.\,data\left(\right)\,,\,\,m\right)\,,\,\,ntt\left(Y.\,data\left(\right)\,,\,\,m\right)\,;
                                                                            fi(0, m) X[i] = X[i] * Y
ntt(X.data(), m, true);
                                                                                                       * Y[i] % P;
void Transform(vector<int> &v, int n) {
                                                                             return X.isz(n() + rhs.n() - 1);
  BitReverse(v, n);
for (int s = 2; s \le n; s \le 1) {
                                                                          Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
     int z = s \gg 1;
for (int i = 0; i < n; i += s) {
                                                                             if (n() = 1) return \{ntt.minv((*this)[0])\};
                                                                             int m = 1;
       for (int k = 0; k < z; ++k) {
                                                                             while (m < n() * 2) m <<= 1;
                                                                             Poly Xi = Poly(*this, (n() + 1) / 2). Inv(). isz(m);
          int x = 1LL
                * v[i + k + z] * omega[kN / s * k] % kMod;
                                                                             Poly Y(*this, m);
          v[i + k + z] = (v[i + k] + kMod' - x) \% kMod;
(v[i + k] += x) \% = kMod;
                                                                             ntt\left(Xi.data\left(\right),\ m\right),\ ntt\left(Y.data\left(\right),\ m\right);
                                                                             if ((Xi[i] \% P) < 0) Xi[i] += P;
  }
                                                                             ntt(Xi.data(), m, true);
void InverseTransform(vector<int> &v, int n) {
                                                                             return Xi.isz(n());
  Transform(v, n);
   for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
                                                                          Poly Sqrt()
   const int kInv = fpow(n, kMod - 2);
                                                                                const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
   for (int i
                                                                             if (n()
         = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;
                                                                                  = 1) return {QuadraticResidue((*this)[0], P)};
```

```
X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
           X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
Poly X(rhs); X. irev(). isz(m);
   Poly Y(*this); Y. irev(). isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m), irev(); X = \text{rhs.Mul}(Q), Y = *\text{this}; fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
   return {Q, Y. isz (max(1, rhs.n() - 1))};
Poly Dx() const {
   Poly ret(n() - 1);
   fi(0,
        ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
   return ret.isz(max(1, ret.n()));
Poly Sx() const {
   Poly ret(n() + 1);
   fi(0, n())
         ret[i + 1] = ntt.minv(i + 1) * (*this)[i] \% P;
   return ret;
Poly _tmul(int nn, const Poly &rhs) const {
   \begin{array}{l} \mbox{Poly } Y = \mbox{Mul}(\mbox{rhs}) . \mbox{ isz} (\mbox{n}() + \mbox{nn} - 1); \\ \mbox{return } \mbox{Poly}(Y. \mbox{data}() + \mbox{n}() - 1, Y. \mbox{data}() + Y. \mbox{n}()); \end{array}
vector<ll> _eval(const
        vector<ll> &x, const vector<Poly> &up) const {
   const int m = (int)x.size();
   if (!m) return { };
   vector Poly> down (m * 2);
   // \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};
   // fi(2, m *
           2) down[i] = down[i / 2].DivMod(up[i]).second;
   down[1] = Poly(up[1])
   \begin{array}{l} . \  \, \mathrm{irev}\,()\,.\,\mathrm{isz}\,(n())\,.\,\mathrm{Inv}\,()\,.\,\mathrm{irev}\,()\,.\,\underline{}\,\mathrm{tmul}(m,\ ^*\mathrm{this})\,;\\ \mathrm{fi}\,(\,2\,,\,m\,^*\,\,2\,)\,\,\mathrm{down}\,[\,i\,]\\ =\,\mathrm{up}\,[\,i\,\,\widehat{}\,\,1\,]\,.\,\underline{}\,\mathrm{tmul}\,(\mathrm{up}\,[\,i\,]\,.\,n()\,\,-\,\,1\,,\,\,\mathrm{down}\,[\,i\,\,/\,\,2])\,; \end{array}
   vectorvectory(m);
   fi(0, m) y[i] = down[m + i][0];
   return y;
static vector<Poly> _treel(const vector<ll> &x) {
  const int m = (int)x.size();
   return up;
}
      <ll> Eval(const vector<ll> &x) const \{ // 1e5, 1s
   auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate (const vector
      <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
   const int m = (int)x.size();

vector<Poly> up = _tree1(x), down(m * 2);

vector<ll> z = up[1].Dx()._eval(x, up);

fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;

fi(0, m) down[m + i] = {z[i]};
   for (int i = m -
         `1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
   return down[1];
Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
   return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] = 0, 1e5/360ms
   if (n() = 1) return \{1\};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
   Poly Y = X.Ln(); Y[0] = P - 1;
   \begin{array}{c} \text{fi} \, (0 \,, \, n()) \\ \text{if} \, ((Y[\,i\,] = (*\,t\,h\,i\,s\,)\,[\,i\,] \, - \, Y[\,i\,]) \, < \, 0) \, \, Y[\,i\,] \, +\!\!= P; \end{array}
   return X. Mul(Y). isz(n());
  / M := P(P - 1). If k >= M, k := k \% M + M.
Poly Pow(ll k) const {
   int nz = 0;
```

```
while (nz < n() \&\& !(*this)[nz]) ++nz; if (nz * min(k, (ll)n()) >= n()) return Poly(n()); if (!k) return Poly(Poly \{1\}, n()); Poly X(data() + nz, data() + nz + n() - nz * k); const ll c = ntt.mpow(X[0], k % (P - 1));
       return X.Ln().imul
             (k \ \% \ P) \ . Exp() \ . imul(c) \ . irev() \ . isz(n()) \ . irev();
   static 11
          LinearRecursion(const vector<ll> &a, const vector
          <11> &coef, ll n) { // a_n = \sum a_n c_j a_n = \sum a_n c_j a_n }
       const int k = (int)a.size();
       assert((int)coef.size() = k + 1);
      Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\}; fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
      C[k] = 1;
      while (n) {
          if (n \% 2) W = W.Mul(M).DivMod(C).second;
          n \neq 2, M = M.Mul(M).DivMod(C).second;
       11 \text{ ret} = 0;
      \mbox{fi} \, (\, 0 \, , \, \, k) \  \  \, \dot{ret} \, = \, (\, ret \, + W[\, i \, ] \  \  \, ^* \, \, a \, [\, i \, ] ) \, \, \% \, \, P; \, \,
      return ret;
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template \Leftrightarrow decltype(Poly\_t::ntt) Poly\_t::ntt = {};
```

7.5 Value Polynomial [fad6e7]

```
struct Poly {
   mint base; // f(x) = poly[x - base]
   vector<mint> poly;
Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
   mint get_val(const mint &x) {
      if (x >= base \&\& x < base + SZ(poly))
         return poly[x - base];
      mint rt = 0;
      \label{eq:continuous} \mbox{vector} < \mbox{mint} > \mbox{lmul}(\mbox{SZ}(\mbox{poly}) \,, \ 1) \,, \ \mbox{rmul}(\mbox{SZ}(\mbox{poly}) \,, \ 1) \,;
      for (int i = 1; i < SZ(poly); +i) lmul[i] = lmul[i - 1] * (x - (base + i - 1)); for (int i = SZ(poly) - 2; i >= 0; --i) rmul[i] = rmul[i + 1] * (x - (base + i + 1)); for (int i = 0.1; SZ(poly) - 1);
      for (int i = 0; i < SZ(poly); ++i)
rt += poly[i] * ifac[i] * inegfac
[SZ(poly) - 1 - i] * lmul[i] * rmul[i];
      return rt;
   mint nw = get_val(base + SZ(poly));
      poly.pb(nw);
      for (int i = 1; i < SZ(poly); ++i)
         poly[i] += poly[i - 1];
```

7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for β being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by Q_k the polynomial such that $F(Q_k)=0\pmod{x^{2^k}}$), then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

8 Geometry

8.1 Basic [b068f0]

```
National Taiwan University 8BQube
  double operator ^ (P b) { return x * b.y - y * b.x; }
                                                                       void dfs (
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P rot(double o) {
     double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
  double angle() { return atan2(y, x); }
                                                                           } else {
struct L {
  // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
 P project(P p) { return pa + (pb - pa).unit
    () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs()) * (pb - pa).abs()); }
bool SegmentIntersect (P p1, P p2, P p3, P p4) {
  if (\max(p1.x, p2.x) < \min(p3.x, p4.x) | |
        \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
  if (\max(p1.y, p2.y) < \min(p3.y, p4.y) | |
        \max(\,\mathrm{p3.y},\,\mathrm{p4.y}) < \min(\,\mathrm{p1.y},\,\,\mathrm{p2.y})) \ \text{return false} \,;
  return sign((p3 - p1)) (p4 - p1)) * sign((p3 - p2)) (p4 - p2)) <= 0 && sign((p1 - p3))
                                                                           return r * r * o / 2;
           (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
     (L x, L y) \{ return same(x.a * y.b, x.b * y.a); \}
P Intersect
     (L\ x,\ L\ y)\ \{\ \textbf{return}\ P(\textbf{-}x.b\ *\ y.c\ +\ x.c\ *\ y.b,\ x
     a * y.c - x.c * y.a / (-x.a * y.b + x.b * y.a); }
8.2 KD Tree [36d550]
namespace kdt {
int root, lc[maxn],
      rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
  if (l = r) return -1;
function<br/>
sool(const point &, const point)
        &>> f = [dep](const point &a, const point &b) {
```

```
if (dep \& 1) return a.x < b.x;
    else return a.y < b.y;
  \inf m = (l + r) >> 1;
  nth\_element(p+l\,,\ p+m,\ p+r\,,\ f)\,;
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc [m] = build (1, m, dep + 1);
if (~lc [m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);

yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);

yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds | | q.x > xr[o] + ds | |
       q.y <
           yl[o] - ds \mid \mid q.y > yr[o] + ds) return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
    (a.y - b.y) * 111 * (a.y - b.y);
```

```
const point &q, long long &d, int o, int dep = 0) {
   if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
  if ((dep & 1)
        && q.x < p[o].x \mid | !(dep & 1) & q.y < p[o].y) {
     if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
long long nearest (const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
8.3 Sector Area [ec8913]
  / calc area of sector which include a, b
double SectorArea(Pa, Pb, double r) {
  double o = atan2(a.y, a.x) - atan2(b.y, b.x);

while (o <= 0) o += 2 * pi;

while (o >= 2 * pi) o -= 2 * pi;

o = min(o, 2 * pi - o);
```

8.4 Half Plane Intersection [0954c1]

```
bool jizz (L l1, L l2, L l3) {
   P = Intersect(12, 13);
    \begin{array}{lll} \hline {\bf return} & ((\, {\bf l}\, {\bf 1}\, .\, {\bf p\dot{b}}\, {\bf -}\, {\bf l}\, \dot{\bf 1}\, .\, {\bf p\dot{a}})\, \hat{\  \  } ({\bf p}\, {\bf -}\, {\bf l}\, {\bf 1}\, .\, {\bf pa})\, ) < & {\bf eps} \, ; \\ \end{array} 
 bool cmp(const L &a, const L &b){
   return same(
         a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
 // availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
 vector<P> HPI(vector<L> &ls){
   sort(ls.begin(),ls.end(),cmp);
    \operatorname{vector} < L > \operatorname{pls}(1, \operatorname{ls}[0]);
   for (int i=0; i<(int) ls. size ();++i) if (!
         same(ls[i].o, pls.back().o))pls.push_back(ls[i]);
   deque < int > dq; dq.push_back(0); dq.push_back(1);
 #define meow(a,b,c
      ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
   for (int i=2;i<(int) pls.size();++i){
    meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
      meow(i,dq[0],dq[1])dq.pop\_front();
      dq.push_back(i);
   meow (dq
          .front(),dq.back(),dq[dq.size()-2])dq.pop_back();
   meow(dq.back(),dq[0],dq[1])dq.pop\_front();
   if (dq.size()<3u) return vector
         <P>(); // no solution or solution is not a convex
    vector<P> rt;
   for (int i=0; i<(int)dq. size();++i)rt.push_back
         (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
   return rt;
}
```

8.5 Rotating Sweep Line [b9fa8d]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
   int n=int(ps.size());
   \text{vector} {<} \text{int} {>} \text{ id} \left( n \right), pos \left( n \right);
   vector<pair<int, int>>> line(n*(n-1)/2);
   int m=-1:
   for(int i=0;i< n;++i)for
           ( \  \, \mathbf{int} \  \  \, \mathbf{j} \! = \! \mathbf{i} \! + \! 1; \mathbf{j} \! < \! \mathbf{n}; \! + \! + \! \mathbf{j} \, ) \, \, \mathbf{line} [ \! + \! + \! \mathbf{m} \! ] \! = \! \mathbf{make\_pair} \, ( \, \mathbf{i} \, \, , \, \mathbf{j} \, ) \, ; \  \, + \! + \! \mathbf{m};
   sort(line.begin(), line.end(), [\&](const
            pair<int, int> &a, const pair<int, int> &b)->bool{
          if (ps
                  [a.first].first==ps[a.second].first)return 0;
          if(ps
                  [b.first].first=ps[b.second].first)return 1;
```

8.6 Triangle Center [33473a]

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) /
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.x + b.x) / 2;

double by = (c.y + b.y) / 2;

double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)) / (\sin(a1) * \cos(a2) - \sin(a2) * \cos(a1));
  return Point (ax + r1 * \cos(a1), ay + r1 * \sin(a1));
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b
       , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);
  res.x = (
       la *`a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (
       la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
}
```

8.7 Polygon Center [728c3a]

```
Point BaryCenter(vector<Point> &p, int n) {
   Point res(0, 0);
   double s = 0.0, t;
   for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
   }
   res.x /= (3 * s);
   res.y /= (3 * s);
   return res;
}</pre>
```

8.8 Maximum Triangle [55b8cb]

```
19
      while (fabs(Cross(p[
            \begin{array}{l} {\rm res} \big[ (\,j + 1) \,\% \,\, {\rm chnum} \big] \,\, - \,\, p \big[ {\rm res} \, [\,i\,] \big] \,\, , \,\, p \big[ {\rm res} \, [\,k\,] \big] \,\, - \,\, p \big[ {\rm res} \, [\,i\,] \big] \,) \,\, > \,\, {\rm fabs} \,(\, {\rm Cross} \, (p \, [\,{\rm res} \, [\,j\,] \big] \,\, - \,\, p \, [\,{\rm res} \, [\,i\,] \big] \,\, , \,\, \\ p \big[ {\rm res} \, [\,k\,] \big] \,\, - \,\, p \big[ {\rm res} \, [\,i\,] \big] \,) \,) \,\, j \,\, = \,\, (j \,\, + \,\, 1) \,\,\% \,\, {\rm chnum} \,; \\ \end{array} 
      tmp = fabs (Cross (
            p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
      if (tmp > area) area = tmp;
   return area / 2;
8.9 Point in Polygon [88cf80]
int pip(vector<P> ps, P p) {
   for (int i = 0; i < ps.size(); ++i) {
      int a = i, b = (i + 1) \% ps. size();
     L l(ps[a], ps[b]);
     P q = l.project(p);
      if ((p - q).abs() < eps && l.inside(q)) return 1;
      if (same(ps[
      [a].y, ps[b].y) \&\& same(ps[a].y, p.y)) continue;
if [ps[a].y > ps[b].y) swap(a, b);
      if (ps[a].y <= p.y && p.y <
            \begin{array}{l} ps[b].y & & p.x <= ps[a].x + (ps[b].x - ps[a].x \\ ) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) ++c; \end{array}
   return (c & 1) * 2;
8.10 Circle [b6844a]
struct C {
  Р с;
   double r;
  C(P\ c = P(\,0\,,\ 0\,)\,,\ \text{double}\ r = \,0)\ :\ c(\,c\,)\,,\ r(\,r\,)\ \{\}
vector<P> Intersect (C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
   vector<P> p;
   if (same(a.r + b.r,
          d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
   else if (a.r + b.r > d \& d + a.r >= b.r) {
     double o = acos
            ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
     P i = (b.c - a.c).unit();
     p.\,push\_back(\,a.\,c\,\,+\,\,i\,.\,rot\,(o)\,\,*\,\,a.\,r\,)\,;
     p.push_back(a.c + i.rot(-o) * a.r);
   return p;
double IntersectArea(Ca, Cb) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
   if (d \ge a.r + b.r - eps) return 0;
   if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
   double p = acos
         ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
   double q = acos
  // remove second
       level if to get points for line (defalut: segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
   double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
```

double C = sq(a.x - o.x)

if $(j - 1.0 \le eps \& j > =$

 $d = \max(0., d);$

vector<P> t;
if (d >= -eps)

return t;

// calc area

) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;

-eps) $t.emplace_back(a.x + i * x, a.y + i * y);$

-eps) t.emplace_back(a.x + j * x, a.y + j * y);

intersect by circle with radius r and triangle OAB

double AreaOfCircleTriangle(Pa, Pb, double r) {

bool ina = a.abs() < r, inb = b.abs() < r;

```
auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
     if (inb) return abs(a ^ b) / 2;
     return SectorArea(b, p[0], r) + abs(a \hat{p}[0]) / 2;
  if (inb) return
         SectorArea(p[0], a, r) + abs(p[0] ^ b)
  if (p.size() = 2u) return SectorArea(a, p[0],
        + SectorArea(p[1], b, r) + abs(p[0]
  else return SectorArea(a, b, r);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
     int j = (i + 1) \% 3;
     double o = atan2
     \begin{array}{lll} (ps\,[\,i\,].\,y,\ ps\,[\,i\,].\,x) \ \ -\ atan2\,(ps\,[\,j\,].\,y,\ ps\,[\,j\,].\,x)\,;\\ if\ (o>=\ pi)\ o=o\ -\ 2\ *\ pi\,;\\ if\ (o<=\ -pi)\ o=o\ +\ 2\ *\ pi\,; \end{array}
     ans += AreaOfCircleTriangle
          (ps[i], ps[j], r) * (o>= 0 ? 1 : -1);
  return abs(ans);
}
```

8.11 Tangent of Circles and Points to Circle [477789]

```
vector <L> tangent (C a, C b) {
#define Pij \
  P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); \
  z.emplace\_back(a.c + i, a.c + i + j);
#define deo(I,J)
  double d = (a)
       .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
  P i =
  vector<L> z;
  if ((a.c - b.c).abs() + b.r < a.r) return z
  else if (same((a.c - b.c).abs() + b.r, a.r)) \{ Pij; \}
  else {
    \operatorname{deo}(-,+);
    if (same(d, a.r + b.r)) { Pij; }
else if (d > a.r + b.r) { deo(+,-); }
  return z;
vector <L> tangent (C c, P p) {
  vector<L> z;
  double d = (p - c.c).abs();
  if_{-}(\mathrm{same}(d\,,\ c\,.\,r\,)\,)\ \{
    P i = (p - c.c).rot(pi / 2);
    z.emplace\_back(p,\ p+i);
  else if (d > c.r) 
    \frac{\text{double o}}{\text{double o}} = a\cos(c.r / d);
    P i = (p - c.c).unit
         ()\;,\;\;j\;=\;i\;.\;rot\,(o)\;\;*\;\;c\,.\;r\;,\;\;k\;=\;i\;.\;rot\,(\,\text{-}o)\;\;*\;\;c\,.\;r\;;
    z.emplace\_back(c.c + j, p);
    z.emplace\_back(c.c + k, p);
  }
  return z;
```

8.12 Area of Union of Circles [0590f1]

8.13 Minimun Distance of 2 Polygons [e9c988]

```
\begin{array}{l} //~p,~q~is~convex\\ double~TwoConvexHullMinDist \end{array}
    for (i =
           0; \ i < n; \ +\!\!+\!\! i\,) \ \ \text{if} \left(P[\,i\,]\,.\,y < P[\,Y\!M\!i\!n\!P\,]\,.\,y\right) \ Y\!M\!i\!n\!P = \,i\,;
    for (i =
           0; i < m; ++i) if(Q[i].y > Q[YMaxQ].y) YMaxQ = i;
   P[n] = P[0], Q[m] = Q[0];
for (int i = 0; i < n; ++i) {
       while (tmp = Cross(
               \begin{array}{l} \text{Q[YMaxQ} + 1] & \text{-P[YMinP} + 1], \text{ P[YMinP]} & \text{-P[YMinP} \\ + 1]) & \text{Cross}(\text{Q[YMaxQ]} & \text{-P[YMinP} + 1], \text{ P[YMinP} \\ \end{array} 
        ] - P[YMinP + 1])) YMaxQ = (YMaxQ + 1) % m;
if (tmp < 0) ans = min(ans, PointToSegDist
               (P[YMinP], P[YMinP + 1], Q[YMaxQ]));
         \begin{array}{l} \textbf{else ans} = \min(\textbf{ans}\,,\,\, TwoSegMinDist(P[\ YMinP]\,,\,\, P[YMinP\,+\,\,1]\,,\,\, Q[YMaxQ]\,,\,\, Q[YMaxQ\,+\,\,1]))\,; \end{array} 
       YMinP = (YMinP + 1) \% n;
    return ans;
}
```

8.14 2D Convex Hull [d97646]

```
bool operator<(const P &a, const P &b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator>(const P &a, const P &b) {
  return same(a.x, b.x) ? a.y > b.y : a.x > b.x;
#define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return
    same(a.x, b.x) ? a.y < b.y : a.x < b.x; });</pre>
      (int i = 0; i < ps.size(); ++i) {
     while (p.size() >= 2 \& crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
     while (p.size() > t \& crx(p[p.size()
          2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double
      x) \ \{ \ return \ same(x, \ 0) \ ? \ 0 \ : \ x > 0 \ ? \ 1 \ : \ -1; \ \}
P\ is LL\, (P\ p1\,,\ P\ p2\,,\ P\ q1\,,\ P\ q2\,)\ \{
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
```

```
struct CH {
  int n;
  vector<P> p, u, d;
 CH() {}
 CH(vector < P > ps) : p(ps) {
    n = ps. size();
    rotate(p.begin
         (), min_element(p.begin(), p.end()), p.end());
    auto t = max_element(p.begin(), p.end());
d = vector<P>(p.begin(), next(t));
    u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find (vector <P> &v, P d) {
    int l = 0, r = v.size();
    while (1 + 5 < r) {

int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;

if (v[L] * d > v[R] * d) r = R;
      else l = L;
    int x = 1;
    for (int i = 1 +
          1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
  int findFarest(P v) {
    if (v.y > 0 | v.y = 0 & v.x > 0) return
          ((int)d.size() - 1 + find(u, v)) \% p.size();
    return find(d, v);
 P get(int 1, int r, Pa, Pb) {
    int s = sgn(crx(a, b, p[l \% n]));
    while (l + 1 < r) {
      int m = (l + r) >> 1;
      if (sgn(crx(a, b, p[m \% n])) == s) l = m;
      else r = m;
    return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
  vector < P > getLineIntersect (P a, P b) {
    int X = findFarest((b - a).rot(pi /
    int Y = findFarest ((a - b).rot(pi / 2));
    if (X > Y) swap(X, Y);
    if (sgn
         return {}; // tangent case falls here
  void update_tangent(P q, int i, int &a, int &b) {
    \begin{array}{ll} if & (sgn(crx(q,\ p[a],\ p[i])) > 0) \ a = i \ ; \\ if & (sgn(crx(q,\ p[b],\ p[i])) < 0) \ b = i \ ; \end{array}
  void bs(int 1, int r, Pq, int &a, int &b) {
    if (l == r) return;
    update_tangent(q, 1 % n, a, b):
     int \ s = sgn(crx(q, \ p[1 \ \% \ n] \ , \ p[(l + 1) \ \% \ n])); 
    while (1 + 1 < r) {
      int m = (1 + r) >> 1;
      if (sgn(crx
           else r = m;
    }
    update_tangent(q, r % n, a, b);
  int x = 1;
  \quad \text{for (int } i = l
       +1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
  return x;
int findFarest(P v) {
if (v.y > 0 \mid | v.y = 0 \& v.x > 0) return
       ((int)d.size() - 1 + find(u, v)) \% p.size();
  return find(d, v);
P get(int 1, int r, Pa, Pb) {
  int s = sgn(crx(a, b, p[1 \% n]));
while (1 + 1 < r) {
    int m = (l + r) >> 1;
    if (sgn(crx(a, b, p[m\%n])) == s) l = m;
    else r = m;
  return isLL(a, b, p[1 \% n], p[(1 + 1) \% n]);
vector <P> getIS (P a, P b) {
  int X = findFarest((b - a).spin(pi / 2));
  int Y = findFarest((a - b).spin(pi / 2));
  \quad \text{if} \ (X>Y) \ \operatorname{swap}(X,\ Y) \, ;
```

```
return {};
void update_tangent(P q, int i, int &a, int &b) {
  \begin{array}{ll} \text{if } (\operatorname{sgn}(\operatorname{crx}(q,\ p[a],\ p[i])) > 0) \ a = i\,;\\ \text{if } (\operatorname{sgn}(\operatorname{crx}(q,\ p[b],\ p[i])) < 0) \ b = i\,; \end{array}
void bs(int 1, int r, Pq, int &a, int &b) {
  if (l == r) return;
  update_tangent(q, 1 % n, a, b);
   \begin{array}{l} {\rm int} \ s = sgn(\,crx\,(q,\ p[\,l\ \%\ n]\,,\ p[\,(\,l\ +\ 1)\ \%\ n]\,)\,)\,; \end{array} 
   while (l + 1 < r) {
     int m = (1 + r) >> 1;
     if (sgn
          (\, crx \, (q,\ p \, [m \, \% \ n] \, , \ p \, [\, (m + \, 1) \, \, \% \ n] \, ) \, = \!\!\!\! - \, s \, ) \ l \, = m;
     else r = m;
  update_tangent(q, r % n, a, b);
auto it
        = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
   if (it -> x == p.x) {
    if (it->y > p.y) return 0;
else if (crx(*prev(it), *it, p) < -eps) return 0;
  it = lower_bound
        (u.begin(), u.end(), P(p.x, 1e12), greater < P > ());
   if (it->x = p.x) {
  if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
  a = b = 0;
  int i
       = lower_bound(d.begin(), d.end(), p) - d.begin();
  bs(0, i, p, a, b);
  bs(i\;,\;d.\,size()\;,\;p,\;a,\;b)\,;
  i = lower\_bound(
       u.begin(), u.end(), p, greater<P>()) - u.begin();
  bs((int
        d.size() - 1, (int)d.size() - 1 + i, p, a, b);
  bs((int)d.size()
         -1 + i, (int)d.size() - 1 + u.size(), p, a, b);
};
8.15 3D Convex Hull [clae8f]
```

```
double
     absvol(const P a, const P b, const P c, const P d) {
  return abs(((b-a)^(c-a))^*(d-a))/6;
struct convex3D {
  static const int maxn=1010;
  struct T{
    int a.b.c:
    bool res;
    T()\{\}
    T(int a, int
         b, int c, bool res=1: a(a), b(b), c(c), res(res){}
  int n,m;
  P p [maxn];
  T f [maxn*8];
  int id [maxn][maxn];
  bool on (T &t,P &q) {
    return ((
        p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q, int a, int b) {
    int g=id[a][b];
    if(f[g].res){
      if(on(f[g],p[q]))dfs(q,g);
        id[q][b]=id[a][q]=id[b][a]=m;
        f[m++]=T(b,a,q,1);
    }
  void dfs(int p,int i){
    f[i].res=0;
    meow(p, f[i].b, f[i].a);
```

```
\begin{array}{l} meow(\,p\,,\,f\,[\,\,i\,\,]\,.\,\,c\,\,,\,f\,[\,\,i\,\,]\,.\,\,b)\,\,;\\ meow(\,p\,,\,f\,[\,\,i\,\,]\,.\,\,a\,,\,f\,[\,\,i\,\,]\,.\,\,c)\,\,; \end{array}
                                                                                    if (l = r) return 1e9;
                                                                                    if (r - l = 1) return dist(p[l], p[r]);
                                                                                    int m = (1 + r) >> 1;
   void operator()(){
                                                                                    double d =
     if (n<4)return;
                                                                                          min(closest_pair(l, m), closest_pair(m + 1, r));
     if([\&](){}
                                                                                    vector<int> vec;
                                                                                    \quad \  \  for (int \ i\!=\!1; i\!<\!n;\!+\!+i) \, if (abs
                (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
                                                                                         fabs\,(p\,[m]\,.\,x\ -\ p\,[\,i\,]\,.\,x)\,<\,d\,;\ --\,i\,)\ vec\,.\,push\_back\,(\,i\,)\,;
                                                                                         (int i = m + 1; i \le r \&\&
           }() || [&](){
                                                                                         fabs(p[m].x - p[i].x) < d; ++i) vec.push_back(i);
           for (int [i=2;i<n;++i) if (abs((p[0]-p[i])
                                                                                    sort(vec.begin(), vec.end()
    , [&](int a, int b) { return p[a].y < p[b].y; });</pre>
                  (p[1]-p[i]) > eps) return swap(p[2],p[i]),0;
                                                                                    for (int i = 0; i < vec.size(); ++i) {
  for (int j = i + 1; j < vec.size()
    && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {</pre>
           return 1;
}() || [&](){
           for (int i
                 =3; i < n; ++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                                                                                         d = min(d, dist(p[vec[i]], p[vec[j]]));
                 *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
           }())return;
                                                                                    return d;
     for (int i=0; i<4;++i){
        T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
        if (on(t,p[i]))swap(t.b,t.c);
                                                                                 9
                                                                                       Else
        id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
                                                                                 9.1 Cyclic Ternary Search* [28a883]
        f[m++]=t;
     for(int i=4;i< n;++i)for
                                                                                   * bool pred(int a, int b);
                                                                                 f(0) \sim f(n-1) is a cyclic-shift U-function
           (int j=0; j < m++j) if (f[j]. res & on(f[j], p[i])) {
                                                                                 return idx s.t. pred(x, idx) is false for all x^*/int\ cyc\_tsearch(int\ n, auto\ pred) {
        dfs\left( \,i\,\,,\,j\,\right) ;
        break:
                                                                                    if (n = 1) return 0;
                                                                                         l = 0, r = n; bool rv = pred(1, 0);
     int mm=m; m=0;
                                                                                   while (r - l > 1) {
   int m = (l + r) / 2;
   if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
     for (int i=0; i \le mm + +i) if (f[i].res) f[m++]=f[i];
   bool same(int i, int j){
     return !(absvol(p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].b]),p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol
                                                                                       else l = m;
                                                                                    return pred(1, r % n) ? 1 : r % n;
           (p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
                                                                                 9.2
                                                                                       \mathbf{Mo's}
                                                                                                       Algorithm(With
                                                                                                                                        modification)
   int faces(){
     int r=0;
for(int i=0;i<m,++i){
                                                                                         [5dec12]
        int iden=1;
                                                                                Mo's Algorithm With modification
        for (int j=0; j< i; ++j) if (same(i,j)) iden=0;
                                                                                 Block: N^{2/3}, Complexity: N^{5/3}
        r+=iden;
     }
                                                                                struct Query {
int L, R, LBid, RBid, T;
int
     return r;
                                                                                    Query(int 1, int r, int t):
} tb;
                                                                                      L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
                                                                                    bool operator < (const Query &q) const {
        Minimum Enclosing Circle [7e5b31]
                                                                                      if (LBid != q.LBid) return LBid < q.LBid; if (RBid != q.RBid) return RBid < q.RBid;
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
   double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5; double d = p0 p1;
                                                                                      return T < b.T;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
                                                                                 void solve(vector<Query> query) {
                                                                                    \operatorname{sort}\left(\operatorname{ALL}\left(\operatorname{query}\right)\right);
   return pt(x, y);
                                                                                    int L=0, R=0, T=-1;
                                                                                    for (auto q : query)
                                                                                      while (T < q.T) addTime(L, R, ++T); // TODO while (T > q.T) subTime(L, R, T--); // TODO while (R < q.R) add(arr[++R]); // TODO
circle min_enclosing(vector<pt> &p) {
  random\_shuffle(p.begin()\,,\ p.end());
                                                                                      while (L > q.L) add(arr[-+L]); // TODO while (R > q.R) sub(arr[R--]); // TODO while (L < q.L) sub(arr[L++]); // TODO
   double r = 0.0;
   pt cent;
   for (int i = 0; i < p.size(); ++i) {
     if (norm2(cent - p[i]) <= r) continue;</pre>
                                                                                       // answer query
     cent = p[i];
     r = 0.0;
                                                                                }
     for (int j = 0; j < i; ++j) {
    if (norm2(cent - p[j]) <= r) continue;
    cent = (p[i] + p[j]) / 2;
                                                                                       Mo's Algorithm On Tree [4a7f74]
                                                                                 9.3
        r = norm2(p[j] - cent);
        for (int k = 0; k < j; +++k) {
                                                                                Mo's Algorithm On Tree
           if \ (norm2(cent - p[k]) <= r) \ continue;\\
                                                                                 Preprocess:
           cent = center(p[i], p[j], p[k]);
                                                                                 1) LCA
           r = norm2(p[k] - cent);
                                                                                2) dfs with in[u] = dft++, out[u] = dft++
                                                                                3)
                                                                                    \operatorname{ord}[\operatorname{in}[u]] = \operatorname{ord}[\operatorname{out}[u]] = u
     }
                                                                                    bitset ⟨MAXN⟩ inset
                                                                                 4)
                                                                                struct Query {
int L, R, LBid, lca;
   return circle(cent, sqrt(r));
                                                                                    Query(int u, int v) {
8.17 Closest Pair [7f292a]
                                                                                       int c = LCA(u, v);
                                                                                      \begin{array}{lll} \mbox{if } (c =\!\!\!\!= u \ | \ | \ c =\!\!\!\!= v) \\ q. \, lca =\!\!\!\!\!= -1, \ q. L =\!\!\!\!\!\!= out[c \ \cap \ u \ \cap \ v] \, , \ q. R =\!\!\!\!\!= out[c]; \end{array}
double closest pair(int l, int r) {
   // p should be sorted
```

else if (out[u] < in[v])

increasingly according to the x-coordinates.

```
q.\,lca\,=\,c\,,\;\;q.\,L\,=\,out\,[\,u\,]\,\,,\;\;q.\,R\,=\,in\,[\,v\,]\,;
       else
          q.lca = c, q.L = out[v], q.R = in[u];
       q.Lid = q.L / blk;
    bool operator<(const Query &q) const {
  if (LBid != q.LBid) return LBid < q.LBid;
       return R < q.R;
 void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
        \begin{array}{ll} \textbf{else} & \textbf{add}(\texttt{arr}[\texttt{x}]) \; ; \; \; // \; \textbf{TODO} \end{array} 
       inset[x] = \sim inset[x];
 void solve(vector<Query> query) {
    sort(ALL(query));
    int L = 0, R = 0;
    for (auto q : query) {
       while (R < q.R) flip(ord[++R]);
while (L > q.L) flip(ord[--L]);
while (R > q.R) flip(ord[R--]);
       while (L < q.L) flip (ord[L++]);
       if (~q.lca) add(arr[q.lca]);
       // answer query
       if (~q.lca) sub(arr[q.lca]);
    }
}
```

9.4 Additional Mo's Algorithm Trick

- Mo's Algorithm With Addition Only
 - Sortqueryssameasthenormal Mo's algorithm.
 - For each query [l,r]:
 - If l/blk = r/blk, brute-force.
 - If $l/blk \neq curL/blk$, initialize $curL := (l/blk + 1) \cdot blk$, curR := curL 1
 - If r > curR, increase curR
 - decrease curL to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
 - Require: Changing answer \equiv adding f([l,r],r+1).
 - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
 - Part1: Answer all f([1,r],r+1) first.
 - Part2: Store $curR \to R$ for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
 - $-\ \ Note: You must do the above symmetrically for the left boundaries.$

9.5 Hilbert Curve [ed5979]

9.6 DynamicConvexTrick* [6a6f6d]

```
// only works for integer coordinates!! maintain max
struct Line {
   mutable ll a, b, p;
  bool operator
        <(const Line &rhs) const { return a < rhs.a; }</pre>
  bool operator <(ll x) const { return p < x; }
struct DynamicHull : multiset<Line, less >>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
        lì b) { return a / b - ((a \hat{b}) < 0 \& a \% b); }
  bool isect(iterator x, iterator y) {
  if (y == end()) { x->p = kInf; return 0; }
          ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
     else x->p = Div(y->b - x->b, x->a - y->a);
     \begin{array}{ll} \textbf{return} & \textbf{x->p} >= \textbf{y->p}; \end{array}
  void addline(ll a, ll b) {
  auto z = insert({a, b, 0}), y = z++, x = y;
     while (isect(y, z)) z = erase(z);
     if (x != begin
          () & isect(--x, y)) isect(x, y = erase(y));
```

```
while ((y = x) != begin
           () && (-x)-p >= y-p isect(x, erase(y));
   ll query(ll x) {
     auto l = *lower_bound(x);
      return l.a * x + l.b;
};
        All LCS* [ae68f0]
9.7
void all_lcs(string s, string t) { // 0-base
   vector < int > h(SZ(t));
   iota(ALL(h), 0);
   for (int a = 0; a < SZ(s); ++a) {
      int v = -1;
      for (int c = 0; c < SZ(t); ++c)
        if(s[a] = t[c] | | h[c] < v)
          swap(h[c], v);
      \begin{array}{l} \text{Swap}(\prod_{i \in I}, \dots, i) \\ \text{// LCS}(s[0, a], t[b, c]) = \\ \text{// } c - b + 1 - sum([h[i]] >= b] | i <= c) \end{array}
      // h[i] might become -1 !!
}
       AdaptiveSimpson* [dc2085]
template<typename Func, typename d = double>
 struct Simpson {
   using pdd = pair < d, d>;
   Func f;
   d eval(pdd 1, pdd r, d fm, d eps) {
   pdd m((1.X + r.X) / 2, fm);
   d s = mix(1, r, fm).second;
     auto [flm, sl] = mix(l, m);
auto [fmr, sr] = mix(m, r);
      d \ delta = sl + sr - s;
      if (abs(delta
           ) <= 15 * eps) return sl + sr + delta / 15;
      return eval(1, m, flm, eps / 2) +
        eval(m, r, fmr, eps / 2);
   d eval(d l, d r, d eps) {
      return eval
           (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
    \begin{array}{l} d \ \ eval2 (d \ l \ , \ d \ r \ , \ d \ eps \ , \ \ int \ k = 997) \ \{ \\ d \ h = (r \ - \ l) \ / \ k \ , \ s = \ 0; \\ for \ (int \ i = \ 0; \ i < k; ++i \ , \ l \ += h) \end{array} 
        s \leftarrow eval(l, l + h, eps / k);
      return s;
 };
 template<typename Func>
Simpson < Func > make\_simpson(Func f) \{ return \{f\}; \}
       Simulated Annealing [b14262]
double factor = 100000;
 const int base = 1e9; // remember to run ~ 10 times
 for (int it = 1; it <= 1000000; ++it) {
      // ans:
          answer, nw: current value, rnd(): mt19937 rnd()
      if (\exp(-(nw - ans))
           ) / factor) >= (double)(rnd() % base) / base)
           ans = nw;
      factor *= 0.99995;
9.10 Tree Hash* [e57357]
 ull seed:
 ull shift(ull x) {
  x = x << 13;
   x = x \gg 7;
   x = x << 17;
   return x:
 ull dfs(int u, int f) {
   ull sum = seed;
   for (int i : G[u])
      if (i!= f)
        sum += shift(dfs(i, u));
   return sum;
```

9.11 Binary Search On Fraction [951597]

```
struct Q {
   \begin{array}{l} \text{ll } p, \ q; \\ Q \ go(Q \ b, \ ll \ d) \ \{ \ return \ \{p + b.p*d, \ q + b.q*d\}; \ \} \end{array}
// returns smallest p/q in [lo, hi] such that // pred(p/q) is true, and 0 <= p,q <= N Q frac_bs(11 N) {
   Q lo{0, 1}, hi{1, 0};
if (pred(lo)) return lo;
    assert (pred (hi));
bool dir = 1, L = 1, H = 1;
    for (; L | | H; dir = !dir) {
       ll len = 0, step = 1;
for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
           \begin{array}{l} \text{if } (Q \text{ mid} = \text{hi.go}(\text{lo}\,,\, \text{len}\, + \text{step})\,; \\ \text{mid.p} > N \mid \mid \text{mid.q} > N \mid \mid \text{dir} \, \widehat{\ } \text{pred}(\text{mid})) \end{array}
              t++;
           else len += step;
       swap(lo, hi = hi.go(lo, len));
       (dir ? L : H) = !!len;
    return dir ? hi : lo;
9.12 Bitset LCS [a82d86]
cin >> n >> m;
for (int i = 1, x; i \le n; ++i)
   cin \gg x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
    cin >> x, (g = f) |= p[x];
    f.shiftLeftByOne(), f.set(0);
    ((f = g - f) = g) \&= g;
cout << f.count() << '\n';
```

10 Python

10.1 Misc