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```
1.2 readchar [0754b0]
```

```
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
static char *p = buf, *end = buf;
   if (p == end) end = buf +
         fread_unlocked(buf, 1, bufsize, stdin), p = buf;
  return *p++;
}
1.3 BigIntIO [ea947e]
  _int128 read() {
       _{\text{int}128} \ \ \mathbf{x} = 0, \ \mathbf{f} = 1;
     char ch = getchar();
     while (ch < '0' | | ch > '9') {
    if (ch == '-') f = -1;
          ch = getchar();
     while (ch >= '0' && ch <= '9') {
    x = x * 10 + ch - '0';
          ch = getchar();
     \textcolor{return}{\texttt{return}} \ x \ * \ f;
void print(__int128 x) {
     if (x < 0) {
         putchar('-');
         x = -x;
     if (x > 9) print(x / 10);
     putchar (x \% 10 + \%);
bool cmp(\underline{\phantom{a}}int128 x, \underline{\phantom{a}}int128 y) { return x > y; }
1.4 Black Magic [d566f1]
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace _
                      _gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef
      tree{<} int \;,\; null\_type \;,\; std::less{<} int{>},\; rb\_tree\_tag
       tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int , int> umap;
typedef priority_queue<int> heap;
int main() {
   // rb tree
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order
(0) == 22); assert(*s.find_by_order(1) == 71);
   assert(s.order_of_key
       (22) = 0; assert (s.order_of_key(71) == 1);
   s.erase(22);
  assert(*s.find_by_order
       (0) = 71; \text{ assert (s.order\_of\_key(71) == 0)};
    / mergable heap
  heap a, b; a.join(b);
  // persistant
  rope < char > r[2];
  r[1] = r[0];
  std::string st = "abc";
r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout \ll r[1].substr(0, 2) \ll std::endl;
  return 0:
```

# 1.5 Pragma Optimization [7b330a]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno, unroll-loops")
#pragma GCC target("sse, sse2, sse3, ssse3, sse4")
#pragma GCC target("popent, abm, mmx, avx, arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

#### .6 Bitset [282252]

```
#include<bits/stdc++.h>
using namespace std;

int main () {
   bitset <4> bit;
```

ECC(int \_n): n(\_n), dft()

, ecnt(), necc(), low(n), dfn(n), bln(n), G(n) {} void add\_edge(int u, int v) {}

```
bit.all(); // all bit is true, ret tru;
bit.any(); // any bit is true, ret true
bit.none(); // all bit is false, ret true
bit.count();
                                                                                             G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
                                                                                          void solve() {
                                                                                             is_bridge.resize(ecnt);
for (int i = 0; i < n; ++i)
      bit.count();
      bit.to_string('0', '1');//with parmeter
      bit.reset(); // set all to true
bit.set(); // set all to false
                                                                                                if (!dfn[i]) dfs(i, -1);
      std::bitset <8> b3{0}, b4{42};
                                                                                      }; // ecc_id(i): bln[i]
      std::hash<std::bitset<8>> hash_fn8;
                                                                                       2.3 SCC* [22afe1]
      hash_fn8(b3); hash_fn8(b4);
}
                                                                                       struct SCC { // 0-base
  int n, dft, nscc;
  vector<int> low, dfn, bln, instack, stk;
\mathbf{2}
       Graph
2.1
       BCC Vertex* [ed8308]
                                                                                          vector<vector<int>>> G;
                                                                                          void dfs(int u) {
struct BCC { // 0-base
int n, dft, nbcc;
                                                                                             low[u] = dfn[u] = ++dft;
                                                                                             \begin{array}{l} instack \left[ u \right] = 1, \ stk.pb(u); \\ for \ (int \ v \ : G[u]) \end{array}
   vector<int> low, dfn, bln, stk, is_ap, cir;
   vector<vector<int>>> G, bcc, nG;
                                                                                                if (!dfn[v])
   void make_bcc(int u) {
                                                                                             \begin{array}{l} dfs(v), \ low[u] = min(low[u], \ low[v]); \\ else \ if \ (instack[v] \&\& \ dfn[v] < dfn[u]) \\ low[u] = min(low[u], \ dfn[v]); \\ if \ (low[u] = dfn[u]) \end{array}
      bcc.emplace_back(1, u);
for (; stk.back() != u; stk.pop_back())
  bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
      stk.pop\_back(), bln[u] = nbcc++;
                                                                                                for (; stk.back() != u; stk.pop_back())
                                                                                                   bln[stk
   void dfs(int u, int f) {
                                                                                                          .back()] = nscc, instack[stk.back()] = 0;
      int child = 0;
                                                                                                instack[u] = 0, bln[u] = nscc++, stk.pop\_back();
      low[u] = dfn[u] = ++dft, stk.pb(u);
                                                                                             }
      for (int v : G[u])
         if (!dfn[v]) {
                                                                                          SCC(int _n): n(_n), dft(), nscc
            dfs(v, u), ++child;

low[u] = min(low[u],
                                                                                                (), low(n), dfn(n), bln(n), instack(n), G(n) {}
                                          low[v]);
                                                                                          void add_edge(int u, int v) {
             \begin{array}{ll} & \text{if } (dfn [u] <= low [v]) \\ & \text{is\_ap}[u] = 1, \ bln [u] = nbcc; \\ & \text{make\_bcc}(v), \ bcc.back().pb(u); \\ \end{array} 
                                                                                             G[u].pb(v);
                                                                                          void solve() {
  for (int i = 0; i < n; ++i)</pre>
         else \ if \ (dfn[v] < dfn[u] & v != f)
                                                                                                if (!dfn[i]) dfs(i);
      \begin{array}{l} low[u] = min(low[u], dfn[v]); \\ if \ (f = -1 \&\& \ child < 2) \ is\_ap[u] = 0; \\ if \ (f = -1 \&\& \ child = 0) \ make\_bcc(u); \end{array}
                                                                                       }; // scc_id(i): bln[i]
                                                                                       2.4 2SAT* [e839e5]
   BCC(int _n): n(_n), dft()
                                                                                       struct SAT { // 0-base
         nbcc(), low(n), dfn(n), bln(n), is_ap(n), G(n)  {}
   void add_edge(int u, int v) {
  G[u].pb(v), G[v].pb(u);
                                                                                          int n;
                                                                                          vector<bool> istrue;
                                                                                          SCC scc;
   void solve() {
  for (int i = 0; i < n; ++i)</pre>
                                                                                          SAT(int _n): n(_n), istrue(n + n), scc(n + n)  {}
                                                                                          int rv(int a) {
                                                                                             return a >= n? a - n : a + n;
         if (!dfn[i]) dfs(i, -1);
                                                                                          void add_clause(int a, int b) {
    scc.add_edge(rv(a), b), scc.add_edge(rv(b), a);
   void block_cut_tree() {
      cir.resize(nbcc);
      for (int i = 0; i < n; ++i)
         if (is_ap[i])
                                                                                          bool solve()
           b \ln [i] = nbcc++;
                                                                                             scc.solve();
                                                                                             for (int i = 0; i < n; ++i) {
  if (scc.bln[i] == scc.bln[i + n]) return false;</pre>
      cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
         for (int j : bcc[i])
                                                                                                istrue[i] = scc.bln[i] < scc.bln[i + n];
  if (is_ap[j])
    nG[i].pb(bln[j]), nG[bln[j]].pb(i);
} // up to 2 * n - 2 nodes!! bln[i] for id
                                                                                                istrue[i + n] = !istrue[i];
                                                                                             return true:
                                                                                       };
2.2 Bridge* [Occada]
                                                                                                MinimumMeanCycle* [4be648]
                                                                                       2.5
{\tt struct} \ ECC \ \{ \ // \ 0\text{-base}
   int n, dft, ecnt, necc;
vector<int> low, dfn, bln, is_bridge, stk;
                                                                                       ll road [N] [N]; // input here
                                                                                       struct MinimumMeanCycle {
                                                                                          11 dp[N + 5][N], n;
   vector<vector<pii>>> G;
   void dfs(int u, int f) {
dfn[u] = low[u] = ++dft, stk.pb(u);
                                                                                          pll_solve() {
                                                                                             11 \ a = -1, \ b = -1, \ L = n + 1;
      for (auto [v, e] : G[u])
if (!dfn[v])
                                                                                             for (int i = 2; i \le L; ++i)
                                                                                                for (int k = 0; k < n; ++k)
      dfs(v, e), low[u] = min(low[u], low[v]);

else if (e!= f)

low[u] = min(low[u], dfn[v]);

if (low[u] == dfn[u]) {
                                                                                                   for (int j = 0; j < n; ++j)
                                                                                                      dp[i][j] =
                                                                                             fill[j]
min(dp[i - 1][k] + road[k][j], dp[i][j]);
for (int i = 0; i < n; ++i) {
  if (dp[L][i] >= INF) continue;
         if (f!= -1) is_bridge[f] = 1;
for (; stk.back()!= u; stk.pop_back())
                                                                                                11 ta = 0, tb = 1;
                                                                                                for (int j = 1; j < n; ++j)
  if (dp[j][i] < INF &&
    ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
    ta = dp[L][i] - dp[j][i], tb = L - j;</pre>
            bln[stk.back()] = necc;
         bln[u] = necc++, stk.pop\_back();
     }
```

if (ta = 0) continue;

if (a = -1) | a \* tb > ta \* b) a = ta, b = tb;

```
if (a != -1) {
    ll g = __gcd(a, b);
    return pll(a / g, b / g);
}
    return pll(-1LL, -1LL);
}
void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
}
};</pre>
```

#### 2.6 Virtual Tree\* [80f7cb]

```
vector<int> vG[N];
int top, st[N];
void insert(int u) {
   if (top = -1) return st[++top] = u, void();
   int p = LCA(st[top], u);
  if (p = st[top]) return st[++top] = u, void(); while (top >= 1 && dep[st[top - 1]] >= dep[p])
   vG[st[top - 1]].pb(st[top]), --top;
if (st[top] != p)
     \overrightarrow{vG}[p] \cdot \overrightarrow{pb}(\overrightarrow{st}[\overrightarrow{top}]), --top, \overrightarrow{st}[++top] = p;
   st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top =
   \mathtt{sort}\left(ALL(\,v\,)\;,\right.
     [\&](int a, int b) \{ return dfn[a] < dfn[b]; \});
   for (int i : v) insert(i);
   while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
   // do something
   reset(v[0]);
```

#### 2.7 Maximum Clique Dynst [4a6b3d]

```
struct MaxClique { // fast when N <= 100
  bitset \langle N \rangle G[N], cs[N];
int ans, sol[N], q, cur[N], d[N], n;
void init(int _n) {
    n\,=\,\underline{\phantom{a}}n\,;
    for (int i = 0; i < n; ++i) G[i].reset();
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  \begin{tabular}{ll} void $$ pre\_dfs(vector<int> \&r, int l, bitset< \gg mask) $$ \{$ \end{tabular} \label{table}
    if (l < 4) {
       for (int i : r) d[i] = (G[i] \& mask).count();
       sort (ALL(r)
             , [\&](int x, int y) \{ return d[x] > d[y]; \});
     vector < int > c(SZ(r));
     int lft = \max(\text{ans} - q + 1, 1), rgt = 1, tp = 0;
     cs[1].reset(), cs[2].reset();
     for (int p : r) {
       int k = 1:
       while ((cs[k] \& G[p]).any()) ++k;
       if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
       if (k < lft) r[tp++] = p;
     for (int k = lft; k \ll rgt; ++k)
       for (int p = cs[k]._Find_first
             ()\;;\;\;p<\;N;\;\;p=\;cs\left[\,k\,\right].\,\_Find\_next\left(\,p\,\right)\,)
         r[tp] = p, c[tp] = k, +tp;
     dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
       int>&r, vector<int>&c, int l, bitset<N> mask) {
     while (!r.empty()) {
       int p = r.back();
       r.pop\_back(), mask[p] = 0;
       if (q + c.back() <= ans) return;
       \operatorname{cur}[q++] = p;
       vector<int> nr;
       for (int i : r) if (G[p][i]) nr.pb(i);
```

```
if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
    else if (q > ans) ans = q, copy_n(cur, q, sol);
    c.pop_back(), --q;
}
}
int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
}
};
```

#### 2.8 Minimum Steiner Tree\* [cbf811]

```
struct SteinerTree { // 0-base
  \begin{array}{l} int \ n, \ dst\left[N\right]\left[N\right], \ dp\left[1 << T\right]\left[N\right], \ tdst\left[N\right]; \\ int \ vcst\left[N\right]; \ // \ the \ cost \ of \ vertexs \end{array}
   void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {
       fill_n (dst[i], n, INF);
       dst[i][i] = vcst[i] = 0;
  void chmin(int &x, int val) {
    x = \min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
     for (int k = 0; k < n; ++k)
       for (int i = 0; i < n; ++i)
          for (int j = 0; j < n; ++j)
            chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
     shortest_path();
     int t = SZ(ter), full = (1 << t) - 1; for (int i = 0; i <= full; ++i)
       fill_n (dp[i], n, INF);
     copy_n(vcst, n, dp[0]);
     for (int msk = 1; msk \le full; ++msk) {
       if (!(msk & (msk - 1))) {
          int who = 1
                        _{
m lg}({
m msk});
          for (int i = 0; i < n; ++i)
            dp [msk
                 ][i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; +++i)
          for (int sub = (
               msk - 1) & msk; sub; sub = (sub - 1) & msk)
            chmin (dp [msk] [i],
                dp[sub][i] + dp[msk \cap sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {
          tdst[i] = INF;
          for (int j = 0; j < n; +++j)
            chmin(tdst[i], dp[msk][j] + dst[j][i]);
       copy_n(tdst, n, dp[msk]);
     return *min_element(dp[full], dp[full] + n);
; // O(V 3^T + V^2 2^T)
```

#### 2.9 Dominator Tree\* [e95beb]

```
struct dominator_tree { // 1-base
  vector < int > G[N] , rG[N];
  int n, pa[N] , dfn[N] , id[N] , Time;
  int semi[N] , idom[N] , best[N];
  vector < int > tree [N]; // dominator_tree
  void init (int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
        G[i].clear(), rG[i].clear();
  }
  void add_edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  }
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
        if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  }
  int find(int y, int x) {</pre>
```

```
if (y \le x) return y;
      int tmp = find (pa[y], x);
      \begin{array}{l} \text{if } (\text{semi} [\text{best}[y]] > \text{semi} [\text{best}[pa[y]]]) \\ \text{best}[y] = \text{best}[pa[y]]; \end{array}
      return pa[y] = tmp;
   void tarjan(int root) {
      Time = 0;
      for (int i = 1; i \le n; ++i) {
         dfn[i] = idom[i] = 0;
         tree[i].clear();
best[i] = semi[i] = i;
       dfs(root);
      for (int i = Time; i > 1; --i) {
         \begin{array}{ll} {\bf int} \  \, u \, = \, id \, [\, i \, ] \, ; \end{array}
          for (auto v : rG[u])
            if (v = dfn[v]) {
               find(v, i);
semi[i] = min(semi[i], semi[best[v]]);
          tree\left[\,semi\left[\,i\,\,\right]\,\right].\,pb\left(\,i\,\right)\,;
          for (auto v : tree[pa[i]]) {
            find(v, pa[i]);
idom[v] =
               semi[best[v]] = pa[i] ? pa[i] : best[v];
          tree [pa[i]]. clear();
      for (int i = 2; i \leftarrow Time; ++i)
         if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
tree[id[idom[i]]].pb(id[i]);
  }
};
```

#### 2.10 Minimum Clique Cover\* [5951ca]

```
struct Clique_Cover { // 0-base, O(n2^n)
  int dp[1 << N];
  void init(int _n) {
     n = n, fill_n (dp, 1 << n, 0);
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
  void add_edge(int u, int v) {
     E[u] = 1 << v, E[v] = 1 << u;
  int solve() {
     for (int i = 0; i < n; ++i)

co[1 << i] = E[i] | (1 << i);

co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {
        int t = i & -i;

dp[i] = -dp[i ^ t];

co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)
        co[i] = (co[i] \& i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {
  int sum = 0; // probabilistic
  for (int i = 0; i < (1 << n); ++i)
    sum += (dp[i] *= co[i]);</pre>
        if (sum) return ans;
     }
     return n;
  }
};
```

#### 2.11 NumberofMaximalClique\* [c163d7]

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all [N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
        for (int j = 1; j <= n; ++j) g[i][j] = 0;
  }
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  }
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
```

```
int u = some[d][0];
for (int i = 0; i < sn; ++i) {
    int v = some[d][i];
    if (g[u][v]) continue;
    int tsn = 0, tnn = 0;
    copy_n(all[d], an, all[d + 1]);
    all[d + 1][an] = v;
    for (int j = 0; j < sn; ++j)
        if (g[v][some[d][j]])
            some[d + 1][tsn++] = some[d][j];
        for (int j = 0; j < nn; ++j)
            if (g[v][none[d][j]])
            none[d + 1][tnn++] = none[d][j];
        dfs(d + 1, an + 1, tsn, tnn);
        some[d][i] = 0, none[d][nn++] = v;
    }
}
int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
}
};</pre>
```

#### 3 Data Structure

## 3.1 Discrete Trick

# 3.3 Interval Container\* [78516e]

}

```
remove intervals from a set of disjoint intervals.
 * Will merge the added interval with
       any overlapping intervals in the set when adding.
 st Intervals are [inclusive, exclusive). st_{/}
set<pii>>::
     iterator addInterval(set<pii>% is, int L, int R) {
  if (L == R) return is.end()
  \begin{array}{lll} \textbf{auto} & \textbf{it} = \textbf{is.lower\_bound}(\{L,\ R\})\,, \ \textbf{before} = \textbf{it}\,; \end{array}
  while (it != is.end() & it->X<= R) {
    R = \max(R, it ->Y);
     before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y>= L) {
    L = \min(L, it ->X);
    R = \max(R, it ->Y);
    is.erase(it);
  return is.insert(before, pii(L, R));
void removeInterval(set<pii> is, int L, int R) {
  if (L == R) return;
  auto it = addInterval(is, L, R);
  \begin{array}{ll} \textbf{auto} & r2 \ = \ it \text{->}Y; \end{array}
  if (it->X == L) is .erase(it);
  else (int\&)it ->Y = L;
  if (R != r2) is .emplace(R, r2);
```

#### 3.4 Leftist Tree [bbd228]

```
struct node {
    ll v, data, sz, sum;
    node *1, *r;
    node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
```

```
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
if (a->data < b->data) swap(a, b);
  a \rightarrow r = merge(a \rightarrow r, b);
   if (V(a->r) > V(a->l)) swap(a->r, a->l);
  a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
  a->sum = sum(a->l) + sum(a->r) + a->data;
   return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
   delete tmp;
3.5 Heavy light Decomposition* [babe8a]
```

```
\begin{array}{lll} \textbf{struct} & \textbf{Heavy\_light\_Decomposition} ~\{~//~~ \textbf{1-base} \\ & \textbf{int} ~n, ~u \textbf{link} [N] ~, ~d \textbf{eep} [N] ~, ~m \textbf{xson} [N] ~, ~w [N] ~, ~p \textbf{a} [N] ~; \end{array}
   int t, pl[N], data[N], val[N]; // val: vertex data
   vector<int>G[N];
   void init(int _n) {
      n = _n;
for (int i = 1; i <= n; ++i)
         G[i].clear(), mxson[i] = 0;
   void add_edge(int a, int b) {
      G[a].pb(b), G[b].pb(a);
   void dfs(int u, int f, int d) {
  w[u] = 1, pa[u] = f, deep[u] = d++;
  for (int &i : G[u])
          if (i != f) {
            dfs(i, u, d), w[u] += w[i];
             if (w[mxson[u]] < w[i]) mxson[u] = i;
   }
    \begin{array}{l} \mbox{void } \mbox{cut(int } \mbox{u, int } \mbox{link)} \ \{ \\ \mbox{data[pl[u] = ++t] = val[u], ulink[u] = link;} \end{array} 
      if (!mxson[u]) return;
      cut(mxson[u], link);
for (int i : G[u])
          if (i != pa[u] & i != mxson[u])
            cut(i, i);
   void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
      int ta = ulink[a], tb = ulink[b], res = 0;
while (ta != tb) {
          if (deep
             [ta] > deep[tb]) swap(ta, tb), swap(a, b);
query(pl[tb], pl[b])
         tb = ulink[b = pa[tb]]
      if (pl[a] > pl[b]) swap(a, b);
       // query(pl[a], pl[b])
};
```

#### Centroid Decomposition\* [4eccaf] 3.6

```
struct Cent_Dec { // 1-base
  vector<pll> G[N];
  pll info [N]; // store info. of itself
pll upinfo [N]; // store info. of climbing up
int n, pa [N], layer [N], sz [N], done [N];
ll dis [__lg(N) + 1][N];
void init (int n)
   void init(int _n) {
      n = _n, layer[0] = -1;
      \begin{array}{ll} fill_n (pa+1, n, 0), fill_n (done+1, n, 0); \\ for (int i = 1; i <= n; ++i) G[i].clear(); \end{array}
   void add_edge(int a, int b, int w) {
      G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
   void get_cent(
      int\ u,\ int\ f\ ,\ int\ \&mx,\ int\ \&c\ ,\ int\ num)\ \{
      int mxsz = 0;
      sz[u] = 1;
      for (pll e : G[u])
          if (!done[e.X] && e.X != f) {
    get_cent(e.X, u, mx, c, num);
             sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
```

```
}
if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
  // if required, add self info or climbing info
  dis[layer[org]][u] = d;
  for ( dl e for )
    for (pll e : G[u])
if (!done[e.X] && e.X != f)
         dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    if (sz[e.X] > sz[c])
         \begin{array}{l} lc = cut(e.X, c, num - sz[c]);\\ else \ lc = cut(e.X, c, sz[e.X]); \end{array}
         upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
          a = pa[a], --ly)
       info[a].X += dis[ly][u], ++info[a].Y;
       if (pa[a])
         upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
  îl query(int u) {
    ll rt = 0;
    for (int a = u, ly = layer[a]; a;
          a = pa[a], --ly
       rt += info[a].X + info[a].Y * dis[ly][u];
       if (pa[a])
           upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
};
```

#### 3.7 LiChaoST\* [4a61ec]

```
ll m, k, id;
L(): id(-1) {}
 class LiChao { // maintain max
private:
 int n; vector<L> nodes;
  void insert(int l, int r, int rt, L ln) {
  int m = (l + r) >> 1;
    if (nodes[rt].id == -1)
      return nodes[rt] = ln, void();
    bool at Left = nodes[rt].at(1) < ln.at(1);
    if (nodes[rt].at(m) < ln.at(m))
      atLeft = 1, swap(nodes[rt], ln);
    if (r - l = 1) return;
if (atLeft) insert(l, m, rt << 1, ln);
    else insert (m, r, rt \ll 1 | 1, ln);
  11 query(int 1, int r, int rt, 11 x)
    int m = (1 + r) \gg 1; ll ret = -INF
    if (nodes[rt].id != -1) ret = nodes[rt].at(x);
    if (r - l = 1) return ret;
    if (x
        < m) return max(ret, query(1, m, rt << 1, x));
   return max(ret, query(m, r, rt \ll 1 | 1, x));
public:
 LiChao(int n_) : n(n_), nodes(n * 4)  {}
  void insert(\overline{L} ln) { insert(0, n, 1, ln); }
  ll query(ll x) { return query(0, n, 1, x); }
```

#### 3.8 Link cut tree\* [5f036a]

```
struct Splay { // xor-sum
  static Splay nil;
Splay *ch[2], *f;
```

```
int val, sum, rev, size;
                                                                               Śplay* lca(Splay *x, Splay *y) {
   Splay (int
         _{\text{val}} = 0) : val(_{\text{val}}), sum(_{\text{val}}), rev(0), size(1)
                                                                                 access(x), root_path(y);
   \{ f = ch[0] = ch[1] = &nil; \}
                                                                                  if (y->f = nil) return y;
   bool isr()
                                                                                 return y->f;
   { return f->ch[0] != this && f->ch[1] != this; }
                                                                               void change(Splay *x, int val) {
   int dir()
  { return f->ch[0] = this ? 0 : 1; } void setCh(Splay *c, int d) {
                                                                                 splay(x), x->val = val, x->pull();
     ch[d] = c;
                                                                               int query (Splay *x, Splay *y) {
     if (c != \&nil) c > f = this;
                                                                                 split(x, y);
     pull();
                                                                                 return y->sum;
   void give_tag(int r) {
                                                                               3.9 KDTree [74016d]
     if (r) swap(ch[0], ch[1]), rev = 1;
                                                                               namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  void push() {
  if (ch[0] != &nil) ch[0]->give_tag(rev);
  if (ch[1] != &nil) ch[1]->give_tag(rev);
                                                                                 yl[maxn], yr[maxn];
                                                                               point p[maxn];
     rev = 0;
                                                                               int build(int l, int r, int dep = 0) {
                                                                                  if (l = r) return -1;
   void pull() {
                                                                                 function < bool (const point &, const point &) f =
     // take care of the nil!
                                                                                    [dep](const point &a, const point &b) {
  if (dep & 1) return a.x < b.x;</pre>
     Figure 1. Size + ch[1] - size + 1; sum = ch[0] - size + ch[1] - size + 1; sum = ch[0] - size + ch[1] - size + 1; if (ch[0] != &nil) ch[0] - f = this; if (ch[1] != &nil) ch[1] - f = this;
                                                                                       else return a.y < b.y;
                                                                                 int m = (1 + r) >> 1;
                                                                                 nth\_element(p + l, p + m, p + r, f);
} Śplay::nil;
Splay *nil = &Splay::nil;
                                                                                 xl[m] = xr[m] = p[m].x;

yl[m] = yr[m] = p[m].y;
void rotate (Splay
                         *x) {
                                                                                 lc[m] = build(1, m, dep + 1);
  Splay *p = x -> f;
                                                                                  if (~lc[m]) {
   int d = x->dir();
                                                                                    xl[m] = min(xl[m], xl[lc[m]]);
   if (!p->isr()) p->f->setCh(x, p->dir());
                                                                                     \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \end{array} 
   else x->f = p->f;
  p->setCh(x->ch[!d], d);
                                                                                    yr[m] = max(yr[m], yr[lc[m]]);
  x->setCh(p, !d);
  p->pull(), x->pull();
                                                                                 rc[m] = build(m + 1, r, dep + 1);
void splay(Splay *x) {
  vector<Splay*> splayVec;
  for (Splay *q = x;; q = q->f) {
                                                                                 if (~rc[m]) {
                                                                                    xl[m] = min(xl[m], xl[rc[m]]);
                                                                                    xr[m] = max(xr[m], xr[rc[m]]);
                                                                                    yl[m] = min(yl[m], yl[rc[m]]);

yr[m] = max(yr[m], yr[rc[m]]);
     splayVec.pb(q);
     if (q->isr()) break;
                                                                                 return m;
   reverse(ALL(splayVec));
   for (auto it : splayVec) it->push();
                                                                               bool bound(const point &q, int o, long long d) {
   while (!x->isr()) {
                                                                                 double ds = sqrt(d + 1.0);
     if (x->f->isr()) rotate(x);
                                                                                 else if (x->dir() = x->f->dir())
     rotate(x->f), rotate(x);
else rotate(x), rotate(x);
                                                                                    return false;
                                                                                 return true;
                                                                              long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
        (a.y - b.y) * 111 * (a.y - b.y);
Splay* access(Splay *x) {
  \begin{array}{l} \text{Splay } *q = \text{nil;} \\ \text{for } (; \text{ x != nil; } \text{ x = x->f}) \end{array}
     \operatorname{splay}(x), x - \operatorname{setCh}(q, 1), q = x;
  return q;
                                                                                 void root_path(Splay *x) { access(x), splay(x); }
                                                                                 \begin{array}{l} \text{long long cd} = \text{dist}(p[o], q); \\ \text{if } (\text{cd } != 0) \text{ d} = \min(d, cd); \end{array}
void chroot(Splay *x){
  root_path(x), x->give_tag(1);
                                                                                  if ((dep & 1) & q.x < p[o].x ||
  x->push(), x->pull();
                                                                                    void split (Splay *x, Splay *y) {
                                                                                    if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
  chroot(x), root_path(y);
                                                                                 } else {
                                                                                    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void link (Splay *x, Splay *y) {
  root_path(x), chroot(y);
x->setCh(y, 1);
                                                                               void init(const vector<point> &v) {
void cut(Splay *x, Splay *y) {
                                                                                 for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  split(x, y);
                                                                                 root = build(0, v.size());
   if (y-size != 5) return;
  y->push();
                                                                               long long nearest (const point &q) {
  y->ch[0] = y->ch[0]->f = nil;
                                                                                 long long res = 1e18;
                                                                                 dfs(q, res, root);
Splay* get_root(Splay *x)
                                                                                 return res:
   for (root_path(x); x->ch[0] != nil; x = x->ch[0])
    x->push();
                                                                                 // namespace kdt
   splay(x);
  return x;
                                                                               3.10 Treap [73589f]
bool conn(Splay *x, Splay *y) {
                                                                              struct node -
  return get\_root(x) == get\_root(y);
                                                                                 int data, sz;
```

```
void up() {
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
   if (!a || !b) return a ? a : b;
   if (!a || !b) return a ? a : b;
  if \ (rand() \ \% \ (sz(a) \, + \, sz(b)) < \, sz(a)) \\
    return b->down(), b->l = merge(a, b->l), b->up(), b;
void split (node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), a->up();
  else b = o, split(o->l, a, b->l, k), <math>b->up();
void split2 (node *o, node *&a, node *&b, int k) {
  if (sz(o) \le k) return a = o, b = 0, void();
  o > down();
  if (sz(o->1) + 1 \le k)
    a = o, split2(o->r, a->r, b, k - sz(o->l) - 1);
  else b = o, split2(o->l, a, b->l, k);
  o > up();
node *kth(node *o, int k) {
  if (k \le sz(o->l)) return kth(o->l, k);
  if (k = sz(o->l) + 1) return o;
  return kth(o->r, k - sz(o->l) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
  if (o->data < key)
    else return Rank(o->l, key);
bool erase (node *&o, int k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->l, o->r);
    delete t;
    return 1;
  node *\&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o > up(), 1 : 0;
void insert(node *&o, int k) {
  node *a, *b;
  split (o, a, b, k),
    o = merge(a, merge(new node(k), b));
void interval(node *&o, int l, int r) {
  node *a, *b, *c;
  split2\,(o,\ a,\ b,\ l\ -\ 1)\,,\ split2\,(b,\ b,\ c,\ r)\,;
  // operate
  o = merge(a, merge(b, c));
```

# Flow/Matching

#### Bipartite Matching\* [f07280]

```
\begin{array}{lll} \textbf{struct} & \textbf{Bipartite\_Matching} ~\{~ / / ~\textbf{0-base} \\ & \textbf{int} ~mp[N] ~, ~mq[N] ~, ~dis [N + 1] ~, ~cur [N] ~, ~l ~, ~r~; \\ & \textbf{vector} {<} \textbf{int} {>} ~G[N + 1] ~; \end{array}
   bool dfs(int u) {
       for (int &i = cur[u]; i < SZ(G[u]); ++i) {
          int e = G[u][i];
          if (mq[e] =
                   | | (dis[mq[e]] = dis[u] + 1 & dfs(mq[e]))
              return mp[mq[e] = u] = e, 1;
       return dis[u] = -1, 0;
   bool bfs() {
       queue<int> q;
       fill_n(dis, l + 1, -1);
for (int i = 0; i < l; ++i)
          if (! ~mp[i])
```

```
q.push(i), dis[i] = 0;
  while (!q.empty())
    int u = q.front();
     q.pop();
     for (int e : G[u])
       if (!~dis[mq[e]])
         q.push(mq[e]), dis[mq[e]] = dis[u] + 1;
  return dis[1] != -1;
int matching() {
  int res = 0;
  fill\_n \, (mp, \ l \, , \ -1) \, , \ fill\_n \, (mq, \ r \, , \ l \, ) \, ;
  while (bfs()) {
     fill_n(cur, 1, 0);
     for (int i = 0; i < 1; ++i)
       res += (!~mp[i] && dfs(i));
  return res; // (i, mp[i] != -1)
void add_edge(int s, int t) { G[s].pb(t); }
void init(int _l, int _r) {
  l = _l, r = _r;
for (int i = 0; i <= 1; ++i)
    G[i].clear();
```

#### 4.2 Kuhn Munkres\* [edf909]

```
void init(int _n) {
    n = \underline{n};
     for (int i = 0; i < n; ++i)
       fill_n (w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
  w[a][b] = wei;
  bool Check(int x) {
     if (vl[x] = 1, \sim fl[x])
     return \operatorname{vr}[\operatorname{qu}[\operatorname{qr}++]] = \operatorname{fl}[x] = 1;
while (\sim x) \operatorname{swap}(x, \operatorname{fr}[\operatorname{fl}[x] = \operatorname{pre}[x]]);
     return 0;
  void bfs(int s) {
     fill_n(slk
          , \dot{n}, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
     ql = qr = 0, qu[qr++] = s, vr[s] = 1;
     for (ll d;;) {
       while (ql < qr)
         for (int x = 0, y = qu[ql++]; x < n; ++x)
if (!vl[x] && slk
                 [x] >= (d = hl[x] + hr[y] - w[x][y])) 
               if (pre[x] = y, d) slk[x] = d;
              else if (!Check(x)) return;
       d = INF;
       for (int x = 0; x < n; ++x)
             (!vl[x] \& d > slk[x]) d = slk[x];
       for (int x = 0; x < n; ++x) {
         if (vl[x]) hl[x] += d;
else slk[x] -= d;
if (vr[x]) hr[x] -= d;
       for (int x = 0; x < n; ++x)
          if (!vl[x] && !slk[x] && !Check(x)) return;
     }
  ll solve() {
     fill_n (fl
     hl[i] = *max\_element(w[i], w[i] + n);
     for (int i = 0; i < n; ++i) bfs(i);
     11 \text{ res} = 0;
     for (int i = 0; i < n; ++i) res += w[i][fl[i]];
     return res;
};
```

#### MincostMaxflow\* [47d9d2]

struct MinCostMaxFlow { // 0-base

```
struct Edge {
      ll from, to, cap, flow, cost, rev;
   } *past[N];
    vector<Edge> G[N];
   int inq [N], n, s, t;

11 dis [N], up [N], pot [N];

bool BellmanFord() {
      fill_n(dis, n, INF), fill_n(inq, n, 0);
      queue < int > q;
      auto relax = [&](int u, ll d, ll cap, Edge *e) {
  if (cap > 0 && dis[u] > d) {
    dis[u] = d, up[u] = cap, past[u] = e;
    dis[u] = d, up[u] = cap, past[u] = e;
}
            if (!inq[u]) inq[u] = 1, q.push(u);
      };
      relax(s, 0, INF, 0);
      while (!q.empty())
         int u = q.front();
         q.pop(), inq[u] = 0;
for (auto &e : G[u]) {
            11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
            relax
                  (e.to, d2, min(up[u], e.cap - e.flow), &e);
      return dis[t] != INF;
   }
   void solve(int
      , int _t, \overline{11} &flow, \overline{11} &cost, bool neg = true) { s = \underline{s}, t = \underline{t}, flow = 0, cost = 0;
      if (neg) BellmanFord(), copy_n(dis, n, pot);
      for (; BellmanFord(); copy_n(dis, n, pot)) {
         for (int
         i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
flow += up[t], cost += up[t] * dis[t];
         for (int i = t; past[i]; i = past[i]->from) {
  auto &e = *past[i];
            e.flow += up[t], G[e.to][e.rev].flow -= up[t];
     }
   }
   void init(int _n) {
     n = \underline{n}, fill_n(pot, n, 0);
      for (int i = 0; i < n; ++i) G[i].clear();
   void add_edge(11 a, 11 b, 11 cap, 11 cost) {
   G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
   G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
};
```

# 4.4 Maximum Simple Graph Matching\*

```
struct Matching { // 0-base
 queue<int> q; int n;
vector<int> fa, s, vis, pre, match;
  vector < vector < int >> G;
  int Find(int u)
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
       if (vis[x] = tk) return x;
       vis[x] = tk;
      x = Find(pre[match[x]]);
  void Blossom(int x, int y, int l) {
    for (; Find(x) != 1; x = pre[y]) {
      pre[x] = y, y = match[x];

if (s[y] == 1) q.push(y), s[y] = 0;
       for (int z: \{x, y\}) if (fa[z] = z) fa[z] = l;
  bool Bfs(int r) {
iota(ALL(fa), 0); fill(ALL(s), -1);
    q = queue < int > (); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
       for (int x = q. front(); int u : G[x])
        if (s[u] = -1) {
if (pre[u] = x, s[u] = 1, match[u] = n) {
             for (int a = u, b = x, last;
b!= n; a = last, b = pre[a])
                last =
                    match[b], match[b] = a, match[a] = b;
             return true;
```

```
q.push(match[u]); s[match[u]] = 0;
         } else if (!s[u] \&\& Find(u) != Find(x)) {
             \begin{array}{ll} \text{int } l = LCA(u, x); \\ Blossom(x, u, l); Blossom(u, x, l); \end{array} 
    }
    return false;
  Matching(int _n) : n(\underline{n}), fa(n + 1), s(n + 1), vis
  (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {} void add_edge(int u, int v)
  \{ G[u].pb(v), G[v].pb(u); \}
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
     return ans;
  } // match[x] == n means not matched
4.5 Maximum Weight Matching* [c80005]
\#define REP(i, l, r) for (int i=(l); i<=(r); ++i)
struct WeightGraph { // 1-based
   struct edge { int u, v, w; }; int n, nx;
  vector<int> lab; vector<vector<edge>>> g;
  vector<int> slk, match, st, pa, S, vis;
  g(nx + 1, vector < edge > (nx + 1)), slk(nx + 1), flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
     match = st = pa = S = vis = slk;
    REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, 0\};
  int E(edge e)
  \{ return \ lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; \}
  void update_slk(int u, int x, int &s)
  \{ if (!s \mid | E(g[u][x]) < E(g[s][x]) | s = u; \} 
  void set_slk(int x) {
    slk[x] = 0;
REP(u, 1, n)
        \mbox{if} \ (g[u][x].w > 0 \ \&\& \ st[u] \ != \ x \ \&\& \ S[st[u]] \ = \ 0) \\
         update_slk(u, x, slk[x]);
  void q_push(int x) {
    if (x \le n) q.push(x);
     else for (int y : flo[x]) q_push(y);
  void set_st(int x, int b) {
    st[x] = b;
     if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
    auto it = find (\overline{ALL}(f), xr);
if (auto pr = it - f.begin(); pr % 2 == 1)
       reverse(1 + ALL(f)), it = f.end() - pr;
    auto res = vector(f.begin(), it);
    return f.erase(f.begin(), it), res;
  void set_match(int u, int v) {
    {\rm match}\,[\,u\,] \;=\; g\,\big[\,u\,]\,[\,v\,]\,.\,v\,;
     if (u \le n) return;
     int xr = flo_from[u][g[u][v].u];
    auto &f = flo [u], z = split_flo (f, xr);

REP(i, 0, SZ(z) - 1) set_match(z[i], z[i^1]);
    set_match(xr, v); f.insert(f.end(), ALL(z));
  void augment(int u, int v) {
    for (;;) {
       int xnv = st[match[u]]; set_match(u, v);
       if (!xnv) return;
       set_match(v = xnv, u = st[pa[xnv]]);
  int lca(int u, int v) {
     static int t = 0; ++t;
     for (++t; u | | v; swap(u, v)) if (u) {
       if (vis[u] == t) return u;
       vis[u] = t, u = st[match[u]];
       if (u) u = st[pa[u]];
    return 0;
  void add_blossom(int u, int o, int v) {
    int b = find(n + 1 + ALL(st), 0) - begin(st);
```

lab[b] = 0, S[b] = 0, match[b] = match[o];

```
\begin{array}{l} vector{<}int{>}\ f = \{o\}; \\ for\ (int\ t\ : \{u,\ v\})\ \{ \end{array}
                                                                                                 REP(u\,,\ 0\,,\ n)\ st\,[u]\,=\,u\,,\ flo\,[u]\,.\,clear\,()\,;
                                                                                                 int w_max = 0;
                                                                                                 reverse(1 + ALL(f));
      for (int x = t, y; x != o; x = st[pa[y]])
         f.pb(x), f.pb(y = st[match[x]]), q_push(y);
                                                                                                    w_max = max(w_max, g[u][v].w);
                                                                                                 fill(ALL(lab), w_max);
   flo[b] = f; set\_st(b, b);
   \begin{aligned} & \text{REP}(\mathbf{x}, \ 1, \ n\mathbf{x}) \ g[b][\mathbf{x}].\mathbf{w} = g[\mathbf{x}][b].\mathbf{w} = 0; \\ & \text{fill} \ (\text{ALL}(\text{flo\_from}[b]) \ , \ 0); \end{aligned} 
                                                                                                 int n_matches = 0; ll tot_weight = 0;
                                                                                                  while (matching()) ++n_matches;
   for (int xs : flo[b]) {
                                                                                                 REP(u, 1, n) if (match[u] \&\& match[u] < u)
      \begin{aligned} & \text{REP}(x, 1, nx) \\ & \text{if } (g[b][x].w = 0 \ | \ E(g[xs][x]) < E(g[b][x])) \end{aligned} 
                                                                                                 tot_weight += g[u][match[u]].w;
return make_pair(tot_weight, n_matches);
            g[b][x] = g[xs][x], g[x][b] = g[x][xs];
      REP(x,
                                                                                              void add_edge(int u, int v, int w)
                                                                                             \{g[u][v].w = g[v][u].w = w; \}
         if (flo_from[xs][x]) flo_from[b][x] = xs;
                                                                                          };
   set_slk(b);
                                                                                          4.6 SW-mincut [90bfe6]
void expand_blossom(int b) {
                                                                                           \begin{array}{lll} \textbf{struct SW} \{ \ // \ global \ min \ cut \, , \ O(V^3) \\ \# define \ REP \ for \ (int \ i = 0; \ i < n; \ +\!\!\!+i) \\ \texttt{static const int } MXN = 514, \ INF = 2147483647; \end{array} 
   for (int x : flo[b]) set_st(x, x);
   int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
   for (int x : split_flo(flo[b], xr)) {
                                                                                              \begin{array}{lll} & \text{int} & \text{vst} \left[ \text{MXN} \right] \,, & \text{edge} \left[ \text{MXN} \right] \left[ \text{MXN} \right] \,, & \text{wei} \left[ \text{MXN} \right] \,; \end{array}
      if (xs = -1) { xs = x; continue; } pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
                                                                                              void init(int n)
                                                                                                REP fill_n (edge[i], n, 0);
      slk[xs] = 0, set\_slk(x), q\_push(x), xs = -1;
                                                                                              void addEdge(int u, int v, int w){
   for (int x : flo[b])
                                                                                                 edge[u][v] += w; edge[v][u] += w;
      if (x = xr) S[x] = 1, pa[x] = pa[b];
      else S[x] = -1, set\_slk(x);
                                                                                             int search (int &s, int &t, int n) {
   st[b] = 0;
                                                                                                 fill_n(vst, n, 0), fill_n(wei, n, 0);
                                                                                                 s = t = -1;
bool on_found_edge(const edge &e) {
                                                                                                 int mx, cur;
   \begin{array}{lll} & \text{if (int } u = st[e.u] \;,\; v = st[e.v];\; S[v] = -1) \; \{\\ & \text{int nu} = st[match[v]];\; pa[v] = e.u;\; S[v] = 1; \end{array}
                                                                                                 for (int j = 0; j < n; +++j) {
                                                                                                   mx = -1, cur = 0;

REP \text{ if } (wei[i] > mx) \text{ } cur = i, mx = wei[i];
   \begin{array}{l} \operatorname{slk}\left[v\right] = \operatorname{slk}\left[\operatorname{nu}\right] = \operatorname{S}\left[\operatorname{nu}\right] = 0; \ \operatorname{q\_push}(\operatorname{nu}); \\ \\ \operatorname{else} \ \operatorname{if} \ \left(\operatorname{S}\left[v\right] = 0\right) \end{array}\right\}
                                                                                                    vst[cur] = 1, wei[cur] = -1;
      if (int o = lca(u, v)) add\_blossom(u, o, v);
                                                                                                    s\,=\,t\,;\ t\,=\,cur
      else return augment(u, v), augment(v, u), true;
                                                                                                   REP if (!vst[i]) wei[i] += edge[cur][i];
   return false;
                                                                                                 return mx;
bool matching() {
                                                                                             int solve(int n) {
   fill(ALL(S), -1), fill(ALL(slk), 0);
                                                                                                 int res = INF;
   q = queue < int > ();
                                                                                                 for (int x, y; n > 1; n--){
  res = min(res, search(x, y, n));
  REP edge[i][x] = (edge[x][i] += edge[y][i]);
  REP(x, 1, nx) \text{ if } (st[x] = x \&\& !match[x])
pa[x] = S[x] = 0, q_push(x);
   if (q.empty()) return false;
                                                                                                    REP {
   for (;;) {
                                                                                                       edge[y][i] = edge[n - 1][i];
      while (SZ(q)) {
                                                                                                    edge[i][y] = edge[i][n - 1];
} // edge[y][y] = 0;
         int u = q. front(); q. pop();
if (S[st[u]] == 1) continue;
         REP(v, 1, n)
                                                                                                 return res;
             if (g[u][v].w > 0 \& st[u] != st[v]) {
                                                                                             }
                if (E(g[u][v]) != 0)
                                                                                          } sw;
                   update_slk(u, st[v], slk[st[v]]);
                                                                                                  BoundedFlow*(Dinic*) [4ae8ab]
                                                                                          4.7
                        (on_found_edge(g[u][v])) return true;
            }
                                                                                          {\color{red} \textbf{struct}} \  \, \textbf{BoundedFlow} \, \, \left\{ \  \, // \, \, \textbf{0-base} \right.
                                                                                             struct edge {
      int d = INF;
                                                                                                int to, cap, flow, rev;
     REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
                                                                                             \label{eq:continuous_continuous_continuous} \begin{array}{l} \text{vector} \! < \! \text{edge} \! > G[N] \, ; \\ \text{int } n, \ s, \ t, \ \text{dis} [N] \, , \ \text{cur} [N] \, , \ \text{cnt} [N] \, ; \end{array}
      REP(x, 1, nx)
         if (int
                                                                                              void init(int _n) {
                  \begin{array}{lll} n = \_n; \\ \text{for (int } i = 0; \ i < n + 2; +\!\!\!+\!\!\!i) \end{array}
            d = min(d, E(g[s][x]) / (S[x] + 2));
                 1, n
                                                                                                    G[i].clear(), cnt[i] = 0;
         if (S[st[u]] == 1) lab[u] += d;
         else if (S[st[u]] == 0) {
  if (lab[u] <= d) return false;
                                                                                              void add_edge(int u, int v, int lcap, int rcap) {
                                                                                                \begin{array}{l} \operatorname{cnt}\left[u\right] \mathrel{-=} \operatorname{lcap}\,, \; \operatorname{cnt}\left[v\right] \mathrel{+=} \operatorname{lcap}\,; \\ G[u].\operatorname{pb}(\operatorname{edge}\left\{v,\;\operatorname{rcap}\,,\;\operatorname{lcap}\,,\;\operatorname{SZ}(G[v])\right\})\,; \end{array}
            lab[u] -= d;
                                                                                                G[v].pb(edge\{u, 0, 0, SZ(G[u]) - 1\});
     REP(b, n + 1, nx) if (st[b] = b \&\& S[b] >= 0)

[ab[b] += d * (2 - 4 * S[b]);
                                                                                             void add_edge(int u, int v, int cap) {
  G[u].pb(edge{v, cap, 0, SZ(G[v])});
  G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
      REP(x, 1, nx)
         if (int s = slk[x]; st[x] == x &&
                s \&\& st[s] != x \&\& E(g[s][x]) == 0
             if (on_found_edge(g[s][x])) return true;
                                                                                             int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < SZ(G[u]); ++i) {
     expand_blossom(b);
                                                                                                    edge \&e = G[u][i]
                                                                                                    if (dis[e.to] = dis[u] + 1 \& e.cap! = e.flow) {
   return false;
                                                                                                        int df = dfs(e.to, min(e.cap - e.flow, cap));
                                                                                                       if (df) {
pair<ll, int> solve() {
                                                                                                           e.\,flow \mathrel{+}= df\,,\; G[\,e.\,to\,]\,[\,e.\,rev\,]\,.\,flow\;-=\;df\,;
   fill\left(ALL(match)\;,\;\;0\right);
                                                                                                           return df;
```

```
}
     \operatorname{dis}[\mathbf{u}] = -1;
     return 0;
  bool bfs() {
     fill_n(dis, n + 3, -1);
     queue<int> q;
     q.push(s), dis[s] = 0;
     while (!q.empty()) {
       int u = q.front();
        q.pop();
        for (edge &e : G[u])
          if (!~dis[e.to] & e.flow != e.cap)
            q.push(e.to), dis[e.to] = dis[u] + 1;
     return dis[t] != -1;
  int maxflow(int _s, int _t) {
     s = _s, t = _t; \\ int flow = 0, df;
     while (bfs()) {
        fill_n(cur, n + 3, 0);
        while ((df = dfs(s, INF))) flow += df;
     return flow;
  bool solve() {
     int sum = 0;
     for (int i = 0; i < n; +++i)
        if (cnt[i] > 0)
        \begin{array}{lll} & \text{add\_edge}(n+1,\ i,\ cnt[i])\ ,\ sum\ +=\ cnt[i];\\ & \text{else if } (cnt[i]<0)\ add\_edge(i,\ n+2,\ -cnt[i])\ ; \end{array}
     if (sum != maxflow(n + 1, n + 2)) sum = -1;
     for (int i = 0; i < n; ++i)
if (cnt[i] > 0)
        G[n + 1].pop_back(); G[i].pop_back(); else if (cnt[i] < 0)
          G[i].pop\_back(), G[n + 2].pop\_back();
     return sum != -1;
  int solve(int _s, int _t) {
    add_edge(_t, _s, INF);

if (!solve()) return -1; // invalid flow

int x = G[_t].back().flow;
     return G[_t].pop_back(), G[_s].pop_back(), x;
};
```

#### 4.8 Gomory Hu tree\* [5f2460]

```
MaxFlow Dinic;
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

#### 4.9 Minimum Cost Circulation\* [cb40c6]

```
struct MinCostCirculation { // 0-base
  struct Edge {
    ll from, to, cap, fcap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  ll\ dis\left[N\right],\ inq\left[N\right],\ n;
  void BellmanFord(int s) {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue < \stackrel{\tt int}{} > q;
    auto relax = [&](int u, ll d, Edge *e) {
       if (dis[u] > d) {
         dis[u] = d, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
    relax(s, 0, 0);
    while (!q.empty())
      int u = q.front();
      q.pop(), inq[u] = 0;
       for (auto &e : G[u])
```

```
if (e.cap > e.flow)
                                          relax(e.to, dis[u] + e.cost, \&e);
                }
        void try_edge(Edge &cur) {
                 if (cur.cap > cur.flow) return ++cur.cap, void();
                BellmanFord(cur.to);
                 if (dis[cur.from] + cur.cost < 0) {
                       ++cur.flow, --G[cur.to][cur.rev].flow;
                         for (int
                              \begin{tabular}{lll} $i=cur.from; &past[i]; &i=past[i]->from) & auto &e=*past[i]; &++e.flow, &--G[e.to][e.rev].flow; & auto &e=*past[i]->from &e=*past[i]-
               ++cur.cap;
         void solve(int mxlg) {
                for (int b = mxlg; b >= 0; --b) {
for (int i = 0; i < n; ++i)
                                for (auto &e : G[i])
                                       e.cap *= 2, e.flow *= 2;
                         for (int i = 0; i < n; ++i)
                                 for (auto &e : G[i])
                                        if (e.fcap >> b & 1)
                                                try_edge(e);
               }
        void init(int _n) { n = _n;
for (int i = 0; i < n; ++i) G[i].clear();</pre>
         void add_edge(ll a, ll b, ll cap, ll cost) {
                G[a].pb(Edge
                                  \{a, b, 0, cap, 0, cost, SZ(G[b]) + (a = b)\}\);
                G[b].pb(Edge\{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1\});
} mcmf; // O(VE * ElogC)
```

# 4.10 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M, x \rightarrow y$  otherwise.
- 2. DFS from unmatched vertices in X.
- 3.  $x \in X$  is chosen iff x is unvisited.
- 4.  $y \in Y$  is chosen iff y is visited.
- Minimum cost cyclic flow
  - 1. Consruct super source S and  $\sinh T$
  - 2. For each edge (x,y,c), connect  $x\to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v) > 0, connect  $S \rightarrow v$  with (cost, cap) = (0, d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - $2. \ \ {\bf Construct\,a\,max\,flow\,model}, {\bf let}\,K\, {\bf be\,the\,sum\,of\,all\,weights}$
  - 3. Connect source  $s \to v, v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G'.

- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
- $3. \ \ The {\it mincut} is equivalent to the {\it maximum} profit of a subset of projects.$
- Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$  , Flow  $f_{uv}$  , Cost  $w_{uv}$  , Required Flow difference for vertex  $b_u$  .
  - 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

```
\begin{aligned} & \min \sum_{uv} w_{uv} f_{uv} \\ & -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ & v_{u} - \sum_{u} f_{uv} = -b_{u} \end{aligned}
```

# 5 String

#### 5.1 KMP [9e1cd1]

```
int F[MAXN];
vector<int> match(string A, string B) {
  vector<int> ans;
  F[0] = -1, F[1] = 0;
  for (int i = 1, j = 0; i < SZ(B); F[++i] = ++j) {
    if (B[i] == B[j]) F[i] = F[j]; // optimize
    while (j != -1 && B[i] != B[j]) j = F[j];
  }
  for (int i = 0, j = 0; i < SZ(A); ++i) {
    while (j != -1 && A[i] != B[j]) j = F[j];
    if (++j == SZ(B)) ans.pb(i + 1 - j), j = F[j];
  }
  return ans;
}</pre>
```

# 5.2 **Z-value\*** [e2dc6f]

```
int z[MAXn]; void make_z(const string &s) { int l = 0, r = 0; for (int i = 1; i < SZ(s); ++i) { for (z[i] = max(0, min(r - i + 1, z[i - l])); i + z[i] < SZ(s) && s[i + z[i]] == s[z[i]]; ++z[i]) ; if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1; } }
```

#### 5.3 Manacher\* [bfe74e]

#### **5.4 SAIS\*** [e9a275]

```
 \begin{vmatrix} \text{auto sais}(\text{const auto \&s}) \ \{ \\ \text{const int } n = SZ(s) \,, \ z = \text{ranges}::\max(s) + 1; \\ \text{if } (n = 1) \ \text{return vector}\{0\}; \\ \text{vector}(\text{int}) \ c(z); \ \text{for } (\text{int } x : s) + + c[x]; \\ \text{partial\_sum}(\text{ALL}(c), \ \text{begin}(c)); \\ \text{vector}(\text{sint}) \ \text{sa}(n); \ \text{auto } I = \text{views}:: \text{iota}(0, n); \\ \text{vector}(\text{sool}) \ t(n, \ \text{true}); \\ \text{for } (\text{int } i = n - 2; \ i >= 0; \ --i) \\ \text{t}[i] = ( \\ \text{s}[i] = s[i + 1] \ ? \ t[i + 1] : \ s[i] < s[i + 1]); \\ \text{auto } \text{is\_lms} = \text{views}:: \text{filter}([\&t](\text{int } x) \ \{ \\ \text{return } x \& t[x] \& ! \ t[x - 1]; \\ \}); \\ \text{auto } \text{induce} = [\&] \ \{ \\ \text{for } (\text{auto } x = c; \ \text{int } y : \text{sa}) \\ \text{if } (y - \cdot) \ \text{if } (! \ t[y]) \ \text{sa}[x[s[y] - 1] + +] = y; \\ \text{for } (\text{auto } x = c; \ \text{int } y : \text{sa} \ | \ \text{views}:: \text{reverse}) \\ \end{aligned}
```

```
if (y--) if (t[y]) sa[--x[s[y]]] = y;
   \label{eq:continuous_section} \begin{array}{lll} \text{vector} <& \text{int} > \text{lms}, \text{ q(n)}; \text{ lms.reserve(n)}; \\ \text{for (auto } x = c; \text{ int i } : \text{ I } | \text{ is\_lms}) \end{array}
       q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
   if (j >= 0) {
          int len = min({n - i, n - j, lms[q[i] + 1] - i});
          ns[q[i]] = nz += lexicographical_compare(
                 \begin{array}{l} \operatorname{begin}(s) + j, & \operatorname{begin}(s) + j + \operatorname{len}, \\ \operatorname{begin}(s) + i, & \operatorname{begin}(s) + i + \operatorname{len}); \end{array}
   fill(ALL(sa), 0); auto nsa = sais(ns);
for (auto x = c; int y : nsa | views::reverse)
      y = lms[y], sa[--x[s[y]]] = y;
   return induce(), sa;
// sa[i]: sa[i]-th suffix
         is the i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
   int n; vector<int> sa, hi, ra;
   Suffix
      (const auto &_s, int _n) : n(_n), hi(n), ra(n) { vector<int> s(n + 1); // s[n] = 0; copy_n(_s, n, begin(s)); // _s shouldn't contain 0
      sa = sais(s); sa.erase(sa.begin());
for (int i = 0; i < n; ++i) ra[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {</pre>
          if (!ra[i]) { h = 0; continue; }
for (int j = sa[ra[i] - 1]; max
                 (i, j) + h < n &  x s[i+h] = s[j+h];) + h;
          hi[ra[i]] = h ? h-- : 0;
};
```

#### 5.5 Aho-Corasick Automatan\* [91c6c0]

```
{\color{red} \textbf{struct}} \hspace{0.1cm} \textbf{AC\_Automatan} \hspace{0.1cm} \{
   int nx[len][sigma], fl[len], cnt[len], ord[len], top;
int rnx[len][sigma]; // node actually be reached
   int newnode() {
      fill_n(nx[top], sigma, -1);
      return top++;
   void init() { top = 1, newnode(); }
   int input(string &s) {
      int X = 1;
      \begin{array}{lll} & \text{for } (c\text{ har } c:s) \ \{ & \text{if } (!\!\sim\!\!\operatorname{nx}[X][c-'A']) \ \operatorname{nx}[X][c-'A'] = \operatorname{newnode}() \,; \\ & X = \operatorname{nx}[X][c-'A'] \,; \end{array}
      return X; // return the end node of string
   void make_fl() {
     queue < int > q;
      q.push(1), fl[1] = 0;
      for (int t = 0; !q.empty(); ) {
         int R = q.front();
        q.pop(), ord [t++] = R;
for (int i = 0; i < sigma; ++i)
            if (~nx[R][i])
              else rnx[R][i] = R > 1 ? rnx[fl[R]][i] : 1;
     }
   void solve() {
      for (int i = top - 2; i > 0; --i)
        cnt[fl[ord[i]]] += cnt[ord[i]];
} ac;
```

#### 5.6 Smallest Rotation [e74dc0]

vector < int > lc(tot);

for (int i = 1; i < tot; ++i) ++lc[len[i]];

```
if (s[i + k] \le s[j + k]) j += k + 1;
                                                                         partial_sum(ALL(lc), lc.begin());
     else i += k + 1;
                                                                         for (int i
    if (i = j) + +j;
                                                                              = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
  int ans = i < n ? i : j;
                                                                      void solve() {
                                                                         for (int i = tot - 2; i >= 0; --i)
  return s.substr(ans, n);
                                                                           cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
5.7 De Bruijn sequence* [f601c2]
                                                                    };
constexpr int MAXC = 10, MAXN = 1e5 + 10;
                                                                    5.9 PalTree* [675736]
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N
                                                                    struct palindromic_tree {
  void dfs(int *out, int t, int p, int &ptr) {
                                                                      struct node {
     if (ptr >= L) return;
                                                                         int next[26], fail, len;
                                                                         int cnt, num; // cnt: appear times, num: number of
    if (t > N) {
       if (N % p) return;
for (int i = 1; i <= p && ptr < L; ++i)
                                                                         // pal. suf.
node(int l = 0): fail(0), len(1), cnt(0), num(0) {
         out[ptr++] = buf[i];
                                                                           for (int i = 0; i < 26; ++i) next[i] = 0;
       else
                                                                         }
       buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
                                                                      };
       for (int j = buf[t - p] + 1; j < C; ++j)
                                                                      vector<node> St;
         buf[t] = j, dfs(out, t + 1, t, ptr);
                                                                      vector < char > s;
                                                                       int last, n;
                                                                       palindromic\_tree() : St(2), last(1), n(0) {
                                                                         St[0].fail = 1, St[1].len = -1, s.pb(-1);
  void solve(int _c, int _n, int _k, int *out) {
    int p = 0;
    \begin{split} C &= \_c, \ N = \_n, \ K = \_k, \ L = N + K - 1; \\ dfs (out, \ 1, \ 1, \ p); \\ if \ (p < L) \ fill (out + p, \ out + L, \ 0); \end{split}
                                                                       inline void clear() {
                                                                         St.clear(), s.clear(), last = 1, n = 0;
                                                                         St.pb(0), St.pb(-1);
                                                                         St[0]. fail = 1, s.pb(-1);
} dbs;
                                                                       inline int get_fail(int x) {
5.8 Extended SAM* [58fa19]
                                                                         while (s[n - St[x].len - 1] != s[n])
                                                                           x = \hat{S}t[x].fail;
struct exSAM { int len [N * 2], link [N * 2]; // maxlength, suflink int next [N * 2] [CNUM], tot; // [0, tot), root = 0 int lenSorted [N * 2]; // topo. order
                                                                         return x;
                                                                      inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
  int cnt[N * 2]; // occurence
                                                                         int cur = get_fail(last);
  int newnode() {
                                                                         if (!St[cur].next[c]) {
    fill_n (next[tot], CNUM, 0);
                                                                           int now = SZ(St);
    len[tot] = cnt[tot] = link[tot] = 0;
                                                                           St.pb(St[cur].len + 2);
    return tot++;
                                                                           St [now]. fail =
                                                                             St[get_fail(St[cur].fail)].next[c];
  void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
                                                                           St[cur].next[c] = now;
                                                                           St[now].num = St[St[now].fail].num + 1;
     int cur = next[last][c];
     len[cur] = len[last] + 1;
                                                                         last = St[cur].next[c], ++St[last].cnt;
    int p = link[last];
while (p != -1 && !next[p][c])
next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
                                                                       inline void count() { // counting cnt
                                                                         auto i = St.rbegin();
                                                                         for (; i != St.rend(); ++i) {
         q = next[p][c];
                                                                           St[i->fail].cnt += i->cnt;
     if (len
                                                                         }
         [p] + 1 = len[q]) return link[cur] = q, cur;
     int clone = newnode()
                                                                       inline int size() { // The number of diff. pal.
     return SZ(St) - 2;
       next[
           clone ][i] = len [next[q][i]] ? next[q][i] : 0;
                                                                    };
    \begin{array}{l} len [clone] = len [p] + 1; \\ while (p!= -1 \&\& next[p][c] == q) \end{array}
                                                                    5.10 Main Lorentz [eaf279]
       next[p][c] = clone, p = link[p];
     link[link[cur] = clone] = link[q];
                                                                    vector<pair<int, int>>> rep[kN]; // 0-base [l, r]
    link[q] = clone;
                                                                    void main_lorentz(const string &s, int sft = 0) {
                                                                       const int n = s.size();
    return cur:
                                                                       if (n = 1) return;
                                                                        const int nu = n / 2, nv = n - nu; 
  void insert(const string &s) {
                                                                      const string u = s.substr(0, nu), v = s.substr(nu),
     int cur = 0;
     for (auto ch : s) {
                                                                             ru(u.rbegin
       int &nxt = next[cur][int(ch - 'a')];
                                                                                   (), u.rend()), rv(v.rbegin(), v.rend());
       \quad \text{if} \quad (\,!\, nxt\,) \quad nxt \,=\, newnode\,(\,)\;;
                                                                      main\_lorentz(u, sft), main\_lorentz(v, sft + nu);
                                                                      const auto z1 = Zalgo(ru), z2 = Zalgo(v + '#' + u), z3 = Zalgo(ru + '#' + rv), z4 = Zalgo(v); auto get_z = [](const vector<int>&z, int i) {
       cnt[cur = nxt] += 1;
    }
  void build() {
                                                                         return
    queue < \stackrel{\textbf{int}}{} > q;
                                                                              (0 \le i \text{ and } i \le (int)z.size()) ? z[i] : 0; };
    q.push(0);
                                                                      auto add_rep
     while (!q.empty())
                                                                            = [&](bool left, int c, int l, int k1, int k2) {
      int cur = q.front();
                                                                               q.pop();
       for (int i = 0; i < CNUM; ++i)
                                                                         if (L > R) return;
                                                                         if (left)
         if (next[cur][i])
           q.push(insertSAM(cur, i));
                                                                               rep[l].emplace\_back(sft + c - R, sft + c - L);
                                                                         else rep[l].emplace_back
```

(sft + c - R - l + 1, sft + c - L - l + 1);

```
for (int cntr = 0; cntr < n; cntr++) {
   int l, k1, k2;
   if (cntr < nu) {
      l = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
   } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
   }
   if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, l, k1, k2);
   }
} // p \in [l, r] \Rightarrow s[p, p + i) = s[p + i, p + 2i)</pre>
```

# 6 Math

## 6.1 ax+by=gcd(only exgcd \*) [5fef50]

```
 \begin{array}{l} pll \ exgcd(ll \ a, \ ll \ b) \ \{ \\ \ if \ (b = 0) \ return \ pll(1, \ 0); \\ \ ll \ p = a \ / \ b; \\ \ pll \ q = exgcd(b, \ a \% \ b); \\ \ return \ pll(q.Y, \ q.X \ - q.Y \ ^* \ p); \\ \} \\ /^* \ ax + by = res \ , \ let \ x \ be \ minimum \ non-negative \\ g, \ p = gcd(a, \ b) \ , \ exgcd(a, \ b) \ ^* \ res \ / \ g \\ \ if \ p.X < 0: \ t = (abs(p.X) + b \ / \ g \ - 1) \ / \ (b \ / \ g) \\ \ else: \ t = -(p.X \ / \ (b \ / \ g)) \\ \ p \ + = \ (b \ / \ g, \ -a \ / \ g) \ ^* \ t \ ^*/ \\ \end{array}
```

#### 6.2 Floor and Ceil [1ffa73]

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 6.3 Floor Enumeration [67ad61]

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
  int x = n / l;
  r = n / x;
}</pre>
```

# **6.4** Mod Min [038fef]

```
// \min\{k \mid 1 \le ((ak) \mod m) \le r\}, no solution -> -1 ll \mod_{\min}(1l \ a, \ ll \ m, \ ll \ l, \ ll \ r) { if (a == 0) return l \ ? -1 : 0; if (1l \ k = (l + a - 1) \ / \ a; \ k * a <= r) return k; ll b = m \ / \ a, \ c = m \% \ a; if (ll \ y = mod_{\min}(c, \ a, \ a - r \% \ a, \ a - l \% \ a)) return (l + y * c + a - 1) \ / \ a + y * b; return -1; }
```

#### 6.5 Linear Mod Inverse [aa1426]

```
 \begin{array}{lll} & & & & & & & & & \\ & & & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ &
```

## 6.6 Linear Filter Mu [663a36]

```
void getMu() {
    mu[1] = 1;
    for (int i = 2; i <= n; ++i) {
        if (!flg[i]) p[++tot] = i, mu[i] = -1;
        for (int j = 1; j <= tot && i * p[j] <= n; ++j) {
            flg[i * p[j]] = 1;
            if (i % p[j] == 0) {
                mu[i * p[j]] = 0;
                break;
            }
            mu[i * p[j]] = -mu[i];
        }
    }
}</pre>
```

#### 6.7 Gaussian integer gcd [4fcbff]

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
     (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
     ll c = a.real() * b.real() + a.imag() * b.imag();
     ll d = a.imag() * b.real() - a.real() * b.imag();
     ll r = b.real() * b.real() + b.imag() * b.imag();
     if (c % r == 0 && d % r == 0) return b;
     return gaussian_gcd
          (b, a - cpx(rnd(c, r), rnd(d, r)) * b);
}
```

#### 6.8 GaussElimination [c016c9]

#### **6.9** Miller Rabin\* [14b81a]

#### 6.10 Simultaneous Equations [21b2e1]

```
struct matrix { //m variables, n equations
   int n, m;
   fraction M[MAXN] [MAXN + 1], sol [MAXN];
   int solve() { //-1: inconsistent, >= 0: rank
for (int i = 0; i < n; ++i) {</pre>
        int piv = 0;
        while (piv < m \& M[i][piv].n) ++piv;
        if (piv == m) continue;
        for (int j = 0; j < n; +++j) {
           if (i = j) continue;
            \begin{array}{ll} \text{fraction } tmp = -M[\,j\,][\,piv\,] \ / \ M[\,i\,][\,piv\,]; \\ \text{for (int } k = 0; \ k <= \end{array} 
                 m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
        }
      int rank = 0;
      for (int i = 0; i < n; ++i) {
        int piv = 0;
        while (piv < m \&\& !M[i][piv].n) ++piv;
        if (piv == m && M[i][m].n) return -1;
        else if (piv
              < m) ++rank, sol[piv] = M[i][m] / M[i][piv];
      return rank:
};
```

### 6.11 Pollard Rho\* [fff0fc]

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
if (prime(n)) return ++cnt[n], void();
  if (n % 2
       = 0) return PollardRho(n / 2), ++cnt[2], void();
  11 x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n) while (true) {
     if (d!= n && d!= 1) {
       PollardRho(n / d);
       PollardRho(d);
       return:
     if (d = n) + p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);

d = gcd(abs(x - y), n);
  }
}
```

#### Simplex Algorithm [40618e]

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a [MAXN] [MAXM] , b [MAXN] , c [MAXM]; double d [MAXN] [MAXM] , x [MAXM]; int ix [MAXN + MAXM]; //!!! array all i
                           !!! array all indexed from 0
// \max\{cx\} \text{ subject to } \{Ax \le b, x > = 0\}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
  value = simplex(a, b, c, N, M);
double simplex (int n, int m) {
  ++m;
  fill_n(d[n], m + 1, 0);
  fill_n(d[n + 1], m + 1, 0);
  iota(ix, ix + n + m, 0);
  int \dot{\mathbf{r}} = \mathbf{n}, \mathbf{s} = \mathbf{m} - 1;
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];
    d[i][m - 1] = 1;

d[i][m] = b[i];
    if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);

d[n + 1][m - 1] = -1;
  for (double dd;; ) {
    if (r < n)
      for (int i = 0; i <= n + 1; ++i) if (i != r) {
  for (int j = 0; j <= m; ++j) if (j != s)
  d[i][j] += d[r][j] * d[i][s];
  d[i][s] *= d[r][s];
      }
    (d[n + 1][j] > -eps & d[n][j] > eps))
           s \; = \; j \; ;
    if(s < 0) break;
    for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {
       if(r < 0 | |
           (dd = d[r][m]
                 (dd < eps \&\& ix[r+m] > ix[i+m]))
    if (r < 0) return -1; // not bounded
  if (d[n + 1][m] < -eps) return -1; // not executable
  double ans = 0:
  fill_n(x, m, 0);
  for (int i = m; i <
    ans += d[i - m][m] * c[ix[i]];
       x[ix[i]] = d[i-m][m];
```

```
return ans:
```

#### 6.12.1 Construction

Primal	Dual
Maximize $c^{\intercal}x$ s.t. $Ax \leq b, x \geq 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \geq c, y \geq 0$
Maximize $c^{\intercal}x$ s.t. $Ax \leq b$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y = c, y \ge 0$
Maximize $c^{\intercal}x$ s.t. $Ax = b, x \ge 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \geq c$

 $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1, n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji} \bar{y}_j = c_i$ holds and for all  $i\!\in\![1,\!m]$  either  $\bar{y}_i\!=\!0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j\!=\!b_j$  holds.

```
1. In case of minimization, let c'_i = -c_i
```

- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### chineseRemainder [fe9f25]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
   11 g = \gcd(m1, m2);
  if ((x^2 - x^1) \% g) return -1; // no sol m1 /= g; m2 /= g;
   \begin{array}{l} pll \ p = exgcd(m1, \ m2) \,; \\ ll \ lcm = m1 \ ^* \ m2 \ ^* \ g \,; \\ ll \ res = p. \ first \ ^* \ (x2 \ - \ x1) \ ^* \ m1 \ + \ x1 \,; \end{array}
   // be careful with overflow
   return (res % lcm + lcm) % lcm;
```

#### 6.14 Factorial without prime factor\* [dcffcb]

```
// O(p^k + log^2 n), pk = p^k
ll prod [MAXP]
ll fac_no_p(ll n, ll p, ll pk) {
   \operatorname{prod}[0] = 1;
   for (int i = 1; i \le pk; ++i)
     if (i \% p) prod[i] = prod[i - 1] * i \% pk;
     else prod[i] = prod[i - 1];
   11 \text{ rt} = 1;
   for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
     rt = rt * prod[n \% pk] \% pk;
   return rt;
} // (n! without factor p) % p^k
```

### 6.15 Discrete Log\* [ba4ac0]

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered\_map{<}int \;,\;\; int{>} \;p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {
    p[\dot{y}] = i;

y = 1LL * y * x % m;
    b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;
    if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m = 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
    if (s == y) return i;
s = 1LL * s * x % m;
  if (s = y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p:
```

#### 6.16 Berlekamp Massey [9380b8]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
    vector \langle T \rangle d(SZ(output) + 1), me, he;
for (int f = 0, i = 1; i <= SZ(output); ++i) {
for (int j = 0; j < SZ(me); ++j)
d[i] += output[i - j - 2] * me[j];
if ((d[i] -= output[i - 1]) == 0) continue;
         if (me.empty()) {
```

```
me. resize(f = i);
   continue;
 vector < T> o(i - f - 1);
T k = -d[i] / d[f]; o.pb(-k);
 return me;
```

#### 6.17**Primes**

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 \ 98789101 \ 987777733 \ 999991921 \ 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 \ \ 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

#### 6.18Theorem

• Cramer's rule

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

• Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i), L_{ij} = -c$  where cisthenumber of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i, j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- Cayley's Formula
  - Given a degree sequence  $d_1, d_2, \dots, d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex 1,2,...,k belong to different components. Then  $T_{n,k}=kn^{n-k-1}.$
- Erdős-Gallaitheorem

As equence of nonnegative integers  $d_1 \ge \cdots \ge d_n$  can be represented as the  ${\tt degree sequence of a finite simple graphon} n vertices if and only if d_1 + \dots + d_n$ 

$$\text{is even and } \sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1 \leq k \leq n.$$

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \dots, b_n$  is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$  holds for

Fulkerson-Chen-Ansteetheorem

A sequence  $(a_1, b_1), \ldots, (a_n, b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \operatorname{holds} \operatorname{for every} 1 \leq k \leq n.$$

• Pick's theorem

For simple polygon, when points are all integer, we have  $A = \#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1$ .

- Möbius inversion formula
  - $\begin{array}{ll} & f(n) = \sum_{d \mid n} g(d) \Leftrightarrow g(n) = \sum_{d \mid n} \mu(d) f(\frac{n}{d}) \\ & f(n) = \sum_{n \mid d} g(d) \Leftrightarrow g(n) = \sum_{n \mid d} \mu(\frac{d}{n}) f(d) \end{array}$
- Sphericalcap
  - A portion of a sphere cut off by a plane.
  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ :
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$  $\cos\theta$ )<sup>2</sup>/3.
  - Area =  $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2 (1 \cos\theta)$ .
- $\bullet \quad Lagrange multiplier$

- Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n) = 0$ .
- Lagrangian function  $\mathcal{L}(x_1, \ldots, x_n, \lambda_1, \ldots, \lambda_k) = f(x_1, \ldots, x_n)$  –
- $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n).$  The solution corresponding to the original constrained optimization
- · Nearest points of two skewlines
  - $\operatorname{Line} 1: \boldsymbol{v}_1 = \boldsymbol{p}_1 + t_1 \boldsymbol{d}_1$
  - Line 2:  $\boldsymbol{v}_2 = \boldsymbol{p}_2 + t_2 \boldsymbol{d}_2$
  - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
  - $\boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n}$

  - $egin{array}{l} & n_1 = a_1 \times n \\ & -n_2 = d_2 \times n \\ & -c_1 = p_1 + rac{(p_2 p_1) \cdot n_2}{d_1 \cdot n_2} d_1 \\ & -c_2 = p_2 + rac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2 \end{array}$
- Derivatives/Integral

Integration by parts: 
$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$

$$\begin{vmatrix} \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}} \\ \int \tan ax = -\frac{\ln|\cos ax|}{a} \end{vmatrix} \frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2}$$

$$\begin{vmatrix} \frac{d}{dx}\tan x = 1 + \tan^2 x \\ \int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \\ \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1) \end{vmatrix}$$

$$\int \sqrt{a^2 + x^2} = \frac{1}{2}\left(x\sqrt{a^2 + x^2} + a^2 \sinh(x/a)\right)$$

$$(x,y,z) = (r\sin\theta\cos\phi, r\sin\theta\sin\phi, r\cos\theta)$$

$$(r,\theta,\phi) = (\sqrt{x^2 + y^2 + z^2}, a\cos(z/\sqrt{x^2 + y^2 + z^2}), a\tan(y,x))$$

• Rotation Matrix

$$M(\theta) = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix}, R_x(\theta_x) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \cos\theta_x & -\sin\theta_x \\ 0 & \sin\theta & \cos\theta \end{bmatrix}$$

#### Estimation 6.19

n |2345678920304050100 p(n) = 23571115223062756044e42e52e8 $n \mid 1001e31e6 1e9 1e12 1e15 1e18$ 

d(i) 12 32 240 1344 6720 26880 103680

n | 1 2 3 4 5 6 7 8 9  $10 \quad \ \ 11 \ \ 12 \ \ 13 \ \ 14 \ \ \ 15$  $\binom{2n}{n}$  2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8  $n \mid 2345678$ 10 11 12 13  $B_n | 2515522038774140211471159757e54e63e7$ 

#### 6.20 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{a} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1)) \\ -h(c,c-b-1,a,m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{c} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \operatorname{mod} c, b \operatorname{mod} c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \operatorname{mod} c, b \operatorname{mod} c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### General Purpose Numbers

Bernoullinumbers  $B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$ 

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into  $\operatorname{exactly} k \operatorname{groups}$ .

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$
 
$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers 1

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. kj:ss.t.  $\pi(j) > \pi(j+1), k+1j$ :ss.t.  $\pi(j) \ge j$ , kj:ss.t.  $\pi(j) > j$ .

$$\begin{array}{l} E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) \\ E(n,0) = E(n,n-1) = 1 \\ E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n} \end{array}$$

# 6.22 Tips for Generating Functions

- Ordinary Generating Function  $A(x) = \sum_{i>0} a_i x^i$ 
  - $-A(rx) \Rightarrow r^n a_n$
  - $-A(x)+B(x) \Rightarrow a_n+b_n$
  - $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} a_i b_{n-i}$
  - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $-xA(x)' \Rightarrow na_n$
  - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function  $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$ 
  - $-A(x)+B(x) \Rightarrow a_n+b_n$

  - $-A^{(k)}(x) \Rightarrow a_{n+k}$   $-A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
  - $-A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $-xA(x) \Rightarrow na_n$
- ${\bf Special \, Generating \, Function}$ 
  - $(1+x)^n = \sum_{i \ge 0} \binom{n}{i} x^i$
  - $rac{1}{(1-x)^n} = \sum_{i \geq 0} \binom{i}{n-1} x^i$   $extbf{Polynomial}$

#### Fast Fourier Transform [9fec80]

```
const int maxn = 131072;
using cplx = complex<double>;
const cplx I = cplx(0, 1);
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
   for (int i = 0; i <= maxn</pre>
         ; ++i) \text{ omega[i]} = \exp(i * 2 * pi / maxn * I);
void bin(vector<cplx> &a, int n) {
    int lg;
    for (lg = 0; (1 << lg) < n; ++lg); --lg;
    vector < cplx > tmp(n);
    for (int i = 0; i < n; ++i) {
         int to = 0;
         for (int j = 0; (1 << j) <
            n; ++j) to |= (((i >> j) \& 1) << (lg - j));
        tmp[to] = a[i];
    for (int i = 0; i < n; ++i) a[i] = tmp[i];
void fft(vector<cplx> &a, int n) {
    bin(a, n);
    for (int step = 2; step \ll n; step \ll 1) {
        int to = step >> 1;
for (int i = 0; i < n; i += step) {
             for (int k = 0; k < to; ++k) {
                 cplx x = a[i
                     + to + k] * omega[maxn / step * k];
                 a[i + to + k] = a[i + k] - x;
```

```
a[i + k] += x;
            }
        }
    }
}
void ifft(vector<cplx> &a, int n) {
    fft(a, n);
    reverse(a.begin() + 1, a.end());
    for (int i = 0; i < n; ++i) a[i] /= n;
vector<int> multiply(const vector<
    int>&a, const vector<int>&b, bool trim = false) {
    int d = 1;
    while
        (d < max(a.size(), b.size())) d <<= 1; d <<= 1;
    vector < cplx > pa(d), pb(d);
    for (int i
         = 0; i < a.size(); ++i) pa[i] = cplx(a[i], 0);
    for (int i
         = 0; i < b.size(); ++i) pb[i] = cplx(b[i], 0);
    fft (pa, d); fft (pb, d);
    for (int i = 0; i < d; ++i) pa[i] *= pb[i];
    ifft (pa, d);
    vector<int> r(d);
    for (int
         if (trim)
        while (r.size() \& r.back() = 0) r.pop_back();
    return r;
                              Root
 Prime
            Root
                  Prime
                  167772161
 7681
            17
                              3
 12289
                  104857601
            11
                              3
 40961
            3
                  985661441
                              3
 65537
            3
                  998244353
                              3
 786433
           10
                  1107296257
                              10
 5767169
                  2013265921
           3
                              31
                  2810183681
 7340033
           3
                              11
 23068673
           3
                  2885681153
                              3
           3
                  605028353
 469762049
```

#### 7.2 Number Theory Transform\* [eleb36]

```
vector<int> omega;
 void Init() {
   omega. resize(kN + 1);
   long long x = \text{fpow}(kRoot, (Mod - 1) / kN);
   omega[0] = 1;
   omega[i] = 1; i <= kN; ++i) {
    omega[i] = 1LL * omega[i - 1] * x % kMod;
void Transform(vector<int> &v, int n) {
   BitReverse(v, n);
for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
     for (int i = 0; i < n; i += s) {
       for (int k = 0; k < z; ++k) {
          v[i + k + z] = (v[i + k] + kMod' - x) \% kMod;
(v[i + k] += x) \% = kMod;
     }
 void InverseTransform(vector<int> &v, int n) {
   Transform \left( v \, , \ n \right);
   for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
   const int kInv = fpow(n, kMod - 2);
   for (int i
         = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;
}
```

#### 7.3 Fast Walsh Transform\* [36c9f5]

```
/* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
   for (int L = 2; L \le n; L \le 1)
      for (int i = 0; i < n; i += L)
for (int j = i; j < i + (L >> 1); ++j)
a[j + (L >> 1)] += a[j] * op;
```

```
National Taiwan University 8BQube
const int N = 21;
int f
     N_{1}^{\dagger}[1 \ll N], g[N][1 \ll N], h[N][1 \ll N], ct[1 \ll N];
                                                                                    Poly Dx() const {
    subset\_convolution(int *a, int *b, int *c, int L) \ \{ \ / \ c_k = \sum_{i=1}^{n} \{i \ | \ j = k, \ i \ \& \ j = 0 \} \ a_i * b_j \ \} \ \}
                                                                                      Poly ret(n() - 1);
                                                                                       fi(0,
   int n = 1 \ll L;
   for (int i = 1; i < n; +++i)
     ct[i] = ct[i \& (i - 1)] + 1;
   for (int i = 0; i < n; ++i)
                                                                                    Poly Sx() const {
     f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
                                                                                      Poly ret(n() + 1);
   for (int i = 0; i \le L; ++i)
                                                                                       fi(0, n())
     fwt(f[i], n, 1), fwt(g[i], n, 1);
   for (int i = 0; i <= L; ++i)
                                                                                      return ret;
     for (int j = 0; j <= i; ++j)

for (int x = 0; x < n; ++x)

h[i][x] += f[j][x] * g[i - j][x];
                                                                                    Polv
   for (int i = 0; i \le L; ++i)
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)
     c[i] = h[ct[i]][i];
                                                                                       if (!m) return { };
7.4 Polynomial Operation [37b8c7]
                                                                                      // fi(2, m *
\begin{array}{ll} fi(s, n) \ for \ (int \ i = (int)(s); \ i < (int)(n); \ +\!\!+\!\!i) \\ template < int \ MAXN, \ ll \ P, \ ll \ RT > /\!/ \ \underline{MAXN} = 2^k \end{array}
struct Poly : vector<ll> { // coefficients in [0, P)
   using vector<ll>>::vector;
   static NTTMAXN, P, RT> ntt;
   int n() const { return (int)size(); } // n() >= 1
                                                                                       vector<ll> y(m);
   Poly(const Poly &p, int m) : vector<ll>(m) {
     copy_n(p.data(), min(p.n(), m), data());
                                                                                      return y;
   Poly& irev()
                                                                                    static vector < Poly>
  { return reverse(data(), data() + n()), *this; } Poly& isz(int m) { return resize(m), *this; }
   Poly& iadd(const Poly &rhs) \{ // n() = rhs.n() \}
     fi(0, n()) if
  (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
return *this;
   Poly& imul(ll k) {
     fi(0, n()) (*this)[i] = (*this)[i] * k % P;
return *this;
                                                                                    vector
   Poly Mul(const Poly &rhs) const {
     int m = 1;
     while (m < n() + rhs.n() - 1) m <<= 1;
     Poly X(*this, m), Y(rhs, m);
     \begin{array}{l} {\rm ntt}\,(X.\,{\rm data}\,()\,,\,\,m)\,,\,\,\,{\rm ntt}\,(Y.\,{\rm data}\,()\,,\,\,m)\,;\\ {\rm fi}\,(\,0\,,\,m)\,\,X[\,{\rm i}\,]\,=\,X[\,{\rm i}\,]\,\,\,{}^*\,\,Y[\,{\rm i}\,]\,\,\%\,\,P; \end{array}
     ntt(X.data(), m, true);
     return X.isz(n() + rhs.n() - 1);
                                                                                       for (int i = m -
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
     if (n() = 1) return \{ntt.minv((*this)[0])\};
                                                                                      return down[1];
     while (m < n() * 2) m <<= 1;
Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
     Poly Y(*this, m);
     ntt(Xi.data(), m), ntt(Y.data(), m);
     fi (0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;
```

ntt(Xi.data(), m, true);

const  $\{ // \text{ Jacobi}((*this)[0], P) = 1, 1e5/235ms \}$ 

= 1) return {QuadraticResidue((\*this)[0], P)};

 $X = Poly(*this \,, \ (n() \,+\, 1) \,\,/\,\, 2).\, Sqrt()\,.\, isz\,(n())\,;$ 

X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);

(const Poly &rhs) const { // (rhs.)back() != 0

if (n() < rhs.n()) return  $\{\{0\}, *this\};$ 

Poly Q = Y.Mul(X.Inv()).isz(m).irev();

const int m = n() - rhs.n() + 1;

Poly X(rhs); X. irev(). isz(m);

Poly Y(\*this); Y. irev(). isz(m);

return Xi.isz(n());

pair<Poly, Poly> DivMod

Poly Sqrt()

**if** (n()

Poly

```
\begin{array}{l} X = rhs.Mul(Q)\,,\; Y = *this\,;\\ fi\,(0\,,\; n())\;\; if\;\; ((Y[\,i\,]\; -=\; X[\,i\,])\,<\,0)\;\; Y[\,i\,]\; +=\, P;\\ return\;\; \{Q,\; Y.\, is\, z\, (max(\,1\,,\; rhs.\, n()\;\; -\; 1))\,\}; \end{array}
         ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
   return ret.isz(max(1, ret.n()));
          ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
        _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn -
   return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const
       vector<ll> &x, const vector<Poly> &up) const {
   const int m = (int)x.size();
   vector < Poly > down (m * 2);
   // \operatorname{down}[1] = \operatorname{DivMod}(\operatorname{up}[1]) \cdot \operatorname{second};
           2) \ \operatorname{down}[\ i\ ] \ = \ \operatorname{down}[\ i\ /\ 2\ ] \ . \ \operatorname{DivMod}(\operatorname{up}[\ i\ ]) \ . \ \operatorname{second};
  down[1] = Poly(up[1])
    .irev().isz(n()).Inv().irev()._tmul(m, *this);
fi(2, m * 2) down[i]
    = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2])
                         1]._tmul(up[i].n() - 1, down[i / 2]);
   fi(0, m) y[i] = down[m + i][0];
                                _{\text{tree1}(\text{const} \text{ vector} < ll > \&x)}  {
   const int m = (int)x.size();
vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = {(x[i] ? P - x[i] : 0), 1};

for (int i = m - 1; i

> 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
     <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
   auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate (const vector
     <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
   const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
fi(0, m) down[m + i] = {z[i]};
         `1; i > 0; --i) down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
Poly Ln() const \{ // (*this)[0] = 1, 1e5/170ms \}
   return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] = 0, 1e5/360ms
  if (n() = 1) return \{1\};

Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());

Poly Y = X.Ln(); Y[0] = P - 1;
   fi(0, n())
           if'((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;
   return X. Mul(Y). isz(n());
 V/M := P(P - 1). If k >= M, k := k \% M + M.
Poly Pow(ll k) const {
   int nz = 0;
   while (nz < n() \&\& !(*this)[nz]) ++nz;
if (nz * min(k, (ll)n()) >= n()) return Poly(n());
if (!k) return Poly(Poly \{1\}, n());
   Poly \ X(data() + nz, \ data() + nz + n() - nz * k);
  const ll c = ntt.mpow(X[0], k % (P - 1));
return X.Ln().imul
         (\texttt{k}~\%~\texttt{P}).\texttt{Exp()}.\texttt{imul(c)}.\texttt{irev()}.\texttt{isz(n())}.\texttt{irev()};\\
static 11
      LinearRecursion(const vector<ll> &a, const vector
      const int k = (int)a.size();
```

```
\begin{array}{l} assert ((int) coef.size() == k + 1); \\ Poly \ C(k + 1), \ W(Poly \ \{1\}, \ k), \ M = \{0, \ 1\}; \\ fi(1, k + 1) \ C[k - i] = coef[i] ? \ P - coef[i] : 0; \\ C[k] = 1; \\ while \ (n) \ \{\\ if \ (n \% \ 2) \ W = W. Mul(M) . DivMod(C) . second; \\ n \ /= 2, \ M = M. Mul(M) . DivMod(C) . second; \\ \} \\ ll \ ret = 0; \\ fi(0, k) \ ret = (ret + W[i] * a[i]) \% \ P; \\ return \ ret; \\ \} \\ \}; \\ \#undef \ fi \\ using \ Poly\_t = Poly < 131072 * 2, 998244353, 3>; \\ template \Leftrightarrow \ decltype(Poly\_t::ntt) \ Poly\_t::ntt = \{\}; \end{array}
```

#### 7.5 Value Polynomial [fad6e7]

```
{\color{red} \textbf{struct}} \hspace{0.1cm} \textbf{Poly} \hspace{0.1cm} \{
  mint base; // f(x) = poly[x - base]
  vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) \{\}
  mint get_val(const mint &x) {
    if (x >= base \&\& x < base + SZ(poly))
      return poly[x - base];
    mint rt = 0;
    for (int i = 0; i < SZ(poly); ++i)
  rt += poly[i] * ifac[i] * inegfac
      [SZ(poly) - 1 - i] * lmul[i] * rmul[i];</pre>
    return rt;
  return;
    mint nw = get_val(base + SZ(poly));
    poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)
       poly[i] += poly[i - 1];
};
```

#### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P) = 0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k) = 0 \pmod{x^{2^k}}$ , then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

# 8 Geometry

#### 8.1 Basic [b068f0]

```
{\color{red} \textbf{double}} \ a, \ b, \ c, \ o;
    P pa, pb;
   \begin{array}{l} P \; pa, \; pb, \\ L() \; : \; a(0), \; b(0), \; c(0), \; o(0), \; pa(), \; pb() \; \{\} \\ L(P \; pa, \; P \; pb) \; : \; a(pa.y \; - \; pb.y), \; b(pb.x \; - \; pa.x \\ \quad \  \  \, ), \; c(pa \; \hat{} \; pb), \; o(atan2(-a, \; b)), \; pa(pa), \; pb(pb) \; \{\} \end{array}
   P project(P p) { return pa + (pb - pa).unit
    () * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
P reflect(P p) { return p + (project(p) - p) * 2; }
double get_ratio(P p) { return (p - pa) * (
    pb - pa) / ((pb - pa).abs()) * (pb - pa).abs()); }
};
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
    if (\max(p1.x, p2.x) < \min(p3.x, p4.x)
             \max(p3.x, p4.x) < \min(p1.x, p2.x)) return false;
    if (\max(p1.y, p2.y) < \min(p3.y, p4.y) | |
             \max(\texttt{p3.y}, \texttt{p4.y}) < \min(\texttt{p1.y}, \texttt{p2.y})) \text{ return false};
   return sign((p3 - p1) ^{\circ} (p4 - p1)) * sign((p3 - p2) ^{\circ} (p4 - p2)) <= 0 && sign((p1 - p3) ^{\circ}
                 (p2 - p3) * sign((p1 - p4) ^ (p2 - p4)) <= 0;
bool parallel
        (L x, L y) { return same(x.a * y.b, x.b * y.a); }
        \begin{array}{l} (L\ x,\ L\ y)\ \{\ {\color{red} return}\ P(\,\hbox{-} x.b\ *\ y.c\ +\ x.c\ *\ y.b,\ x\\ .a\ *\ y.c\ -\ x.c\ *\ y.a)\ /\ (\,\hbox{-} x.a\ *\ y.b\ +\ x.b\ *\ y.a)\,;\ \end{array}\}
8.2 KD Tree [36d550]
namespace kdt {
int root, lc [maxn],
         rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
    if (l = r) return -1;
    function < bool (const point &, const point
            &> f = [dep](const point &a, const point &b) {
        if (dep \& 1) return a.x < b.x;
       else return a.y < b.y;
    int m = (1 + r) >> 1;
    nth\_element(p+1, p+m, p+r, f);
    \begin{array}{l} \operatorname{xl}\left[m\right] = \operatorname{xr}\left[m\right] = \operatorname{p}\left[m\right].\,\mathrm{x}\,;\\ \operatorname{yl}\left[m\right] = \operatorname{yr}\left[m\right] = \operatorname{p}\left[m\right].\,\mathrm{y}\,; \end{array}
    lc[m] = build(1, m, dep + 1);
    if (~lc[m]) {
       xl[m] = min(xl[m], xl[lc[m]]);
        \begin{array}{l} \operatorname{xr}\left[m\right] = \operatorname{max}\left(\operatorname{xr}\left[m\right], \ \operatorname{xr}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \operatorname{yl}\left[m\right] = \operatorname{min}\left(\operatorname{yl}\left[m\right], \ \operatorname{yl}\left[\operatorname{lc}\left[m\right]\right]\right); \\ \end{array} 
       yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
       xl[m] = min(xl[m], xl[rc[m]]);
       \operatorname{xr}[m] = \max(\operatorname{xr}[m], \operatorname{xr}[\operatorname{rc}[m]]);
       yl[m] = min(yl[m], yl[rc[m]]);

yr[m] = max(yr[m], yr[rc[m]]);
    return m:
bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
     if \ (q.x < xl [o] - ds \ | \ | \ q.x > xr [o] + ds \ | \ | \\
                   yl[o] - ds \mid | q.y > yr[o] + ds) return false;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
       const point &q, long long &d, int o, int dep = 0) {
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1)
        && q.x < p[o].x || !(dep & 1) && q.y < p[o].y) { if (\sim lc[o]) dfs(q, d, lc[o], dep + 1);
        if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
    } else {
        if (\sim rc[o]) dfs(q, d, rc[o], dep + 1);
        if (\sim lc [o]) dfs (q, d, lc [o], dep + 1);
```

```
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
}
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
}}
8.3 Sector Area [ec8913]</pre>
```

# // calc area of sector which include a, b double SectorArea(P a, P b, double r) { double o = atan2(a.y, a.x) - atan2(b.y, b.x); while (o <= 0) o += 2 \* pi; while (o >= 2 \* pi) o -= 2 \* pi; o = min(0, 2 \* pi - o); return r \* r \* o / 2;

# 8.4 Half Plane Intersection [0954c1]

```
bool jizz(L l1,L l2,L l3){
  P p=Intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;
bool cmp(const L &a, const L &b){
  return same(
        a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
vector<L> pls(1,ls[0]);
  for (int i=0; i<(int) ls.size();++i) if(!
        same(ls[i].o,pls.back().o))pls.push_back(ls[i]);
  deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c
     ) while (dq. size ()>1u && jizz (pls [a], pls [b], pls [c]))
  for (int i=2; i < (int) pls. size(); ++i)
    \begin{array}{l} meow(\,i\,\,,dq\,.\,back\,(\,)\,\,,dq\,[\,dq\,.\,size\,(\,)\,\,-\,2\,]\,)dq\,.\,pop\_back\,(\,)\,\,;\\ meow(\,i\,\,,dq\,[\,0\,]\,,dq\,[\,1\,]\,)dq\,.\,pop\_front\,(\,)\,\,; \end{array}
    dq.push_back(i);
  }
  meow(dq
        .front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop\_front();\\
  if (dq. size ()<3u) return vector
        <P>(); // no solution or solution is not a convex
  vector<P> rt;
  for (int i=0; i<(int)dq.size();++i)rt.push_back
        (Intersect(pls[dq[i]], pls[dq[(i+1)%dq.size()]]));
  return rt:
}
```

#### 8.5 Rotating Sweep Line [b9fa8d]

```
void rotatingSweepLine(vector<pair<int,int>>> &ps){
  int n=int(ps.size());
  vector < int > id(n), pos(n);
  vector<pair<int, int>>> line(n*(n-1)/2);
  int m=-1:
  for(int i=0;i< n;++i)for
       (int j=i+1; j< n; ++j) line[++m] = make_pair(i,j); ++m;
  sort(line.begin(),line.end(),[&](const
       pair<int, int> &a, const pair<int, int> &b)->bool{
      if (ps
           [a.first].first = ps[a.second].first)return 0;
      if (ps
           [b. first]. first=ps[b. second]. first)return 1;
      return (double
           )(ps[a.first].second-ps[a.second].second)/(ps
           [a. first]. first - ps[a. second]. first) < (double)
           ) (ps[b.first].second-ps[b.second].second
           )/(ps[b.first].first-ps[b.second].first);
      });
  for (int i=0; i< n; ++i) id [i]=i;
  sort(id.begin(),id.end(),[&](const
       int &a, const int &b) { return ps[a] < ps[b]; });
  for (int i=0; i< n; ++i) pos[id[i]]=i;
  for (int i=0; i \leqslant m++i){
    auto l=line[i];
```

```
// meow
    tie (pos[1.first],pos[1.second],
        id [pos[1.first]],id [pos[1.second]])=make_tuple
        (pos[1.second],pos[1.first],1.second,1.first);
}
```

#### 8.6 Triangle Center [33473a]

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y)
  double bx = (c.x + b.x) /
  double by = (c.y + b.y) / 2;
double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by
- ay)) / (sin(a1) * cos(a2) - sin(a2) * cos(a1));
return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point\ TriangleOrthoCenter(Point\ a,\ Point\ b,\ Point\ c)\ \{
  return TriangleMassCenter(a, b
       , c) * 3.0 - TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);
  res.x = (
      la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (
    la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

# 8.7 Polygon Center [728c3a]

```
Point BaryCenter(vector<Point> &p, int n) {
   Point res(0, 0);
   double s = 0.0, t;
   for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
   }
  res.x /= (3 * s);
  res.y /= (3 * s);
  return res;
}</pre>
```

### 8.8 Maximum Triangle [55b8cb]

```
double ConvexHullMaxTriangleArea
                                              (Point p[], int res[], int chnum) {
                         double area = 0, tmp;
                       res[chnum] = res[0];
                       for (int i = 0, j = 1, k = 2; i < chnum; i++) {
                                              while (fabs(Cross(p[
                                                                                      \begin{array}{lll} res[j]] & -p[res[i]] \; , \; p[res[(k+1) \; \% \; chnum]] \; -p[res[i]])) > fabs(Cross(p[res[j]] \; -p[res[i]] \; , \\ p[res[k]] \; -p[res[i]]))) \; k = (k+1) \; \% \; chnum; \\ \end{array} 
                                          tmp = fabs (Cross (
                                                                                        p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                              if (tmp > area) area = tmp;
                                              while (fabs(Cross(p[
                                                                                         \begin{array}{l} \operatorname{res}\left[\left(j+1\right)\%\operatorname{chnum}\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \ \operatorname{p}\left[\operatorname{res}\left[k\right]\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \\ \operatorname{p}\left[\operatorname{res}\left[i\right]\right]\right) > \operatorname{fabs}\left(\operatorname{Cross}\left(\operatorname{p}\left[\operatorname{res}\left[j\right]\right] - \operatorname{p}\left[\operatorname{res}\left[i\right]\right], \\ \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right] - \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right] - \operatorname{chnum}\left[\operatorname{chnum}\left[i\right]\right], \\ \operatorname{chnum}\left[\operatorname
                                                                                               p[res[k]] - p[res[i]]))) j = (j + 1) % chnum;
                                            tmp = fabs (Cross (
                                                                                        p[res[j]] - p[res[i]], p[res[k]] - p[res[i]]));
                                              if (tmp > area) area = tmp;
                      return area / 2;
```

#### 8.9 Point in Polygon [88cf80]

```
int pip(vector<P> ps, P p) {
  int c = 0;
```

for (int i = 0; i < 3; ++i) {

```
for (int i = 0; i < ps. size(); ++i)
                                                                      int j = (i + 1) \% 3;
    int a = i, b = (i + 1) \% ps.size();
                                                                      double o = atan2
                                                                       L l(ps[a], ps[b]);
                                                                      if (o >= pi) o = o - 2 * pi;
if (o <= -pi) o = o + 2 * pi;
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
                                                                      ans += AreaOfCircleTriangle
    if (same(ps[
                                                                           (ps[i]\,,\ ps[j]\,,\ r)\ *\ (o>=0\ ?\ 1\ :\ -1);
         a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
    if (ps[a].y > ps[b].y) swap(a, b);
    if (ps[a].y \le p.y \& p.y \le
                                                                    return abs(ans);
         \begin{array}{l} ps[b].y & & ps[a].x + (ps[b].x - ps[a].x \\ ) / (ps[b].y - ps[a].y) * (p.y - ps[a].y)) + c; \end{array}
                                                                  8.11
                                                                         Tangent of Circles and Points to Circle
  return (c & 1) * 2;
                                                                  vector <L> tangent (C a, C b) {
                                                                  #define Pij \
8.10 Circle [b6844a]
                                                                    P \ i \ = \ (b.c \ - \ a.c) \ . \ unit \ () \ * \ a.r \ , \ j \ = \ P(\,i \ .y \ , \ -i \ .x) \ ; \ \backslash
struct C {
                                                                    z.emplace\_back(a.c + i, a.c + i + j);
 Р с;
                                                                  #define deo(I,J) \
  double r;
                                                                    double d = (a)
  C(P \ c = P(0, 0), \ double \ r = 0) : c(c), \ r(r) \ \{\}
                                                                        .c - b.c).abs(), e = a.r I b.r, o = acos(e / d);
                                                                    vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
                                                                    if (a.r < b.r) swap(a, b);
  vector<P> p;
                                                                    vector<L> z;
  if (same(a.r + b.r,
                                                                    if ((a.c - b.c).abs() + b.r < a.r) return z
        d)) p.push_back(a.c + (b.c - a.c).unit() * a.r);
                                                                    else if (same((a.c - b.c).abs() + b.r, a.r)) \{ Pij; \}
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
                                                                    else {
    double o = acos
                                                                      \operatorname{deo}(-,+);
         \left(\left(sq(a.r) \,+\, sq(d) \,-\, sq(b.r)\right) \,\,/\,\, \left(2 \,\,{}^*\,\, a.r \,\,{}^*\,\, d\right)\right);
                                                                      if (same(d, a.r + b.r)) \{ Pij; \}
else if (d > a.r + b.r) \{ deo(+,-); \}
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.rot(o) * a.r);
    p.push_back(a.c + i.rot(-o) * a.r);
                                                                    return z;
                                                                  }
  return p;
                                                                  vector<L> tangent(C c, P p) {
double IntersectArea(C a, C b) {
                                                                    vector < L > z;
  if (a.r > b.r) swap(a, b);
                                                                    \frac{double \ d = (p - c.c).abs();}{}
  double d = (a.c - b.c).abs();
                                                                    if(same(d, c.r)) {
  if (d \ge a.r + b.r - eps) return 0;
                                                                      P i = (p - c.c).rot(pi / 2);
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
                                                                      z.emplace\_back(p, p + i);
  double p = a\cos
                                                                    else if (d > c.r) 
       ((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
                                                                       double o = acos(c.r / d);
  double q = acos
                                                                      P i = (p - c.c) \cdot unit'
  (), j = i.rot(o) * c.r, k = i.rot(-o) * c.r;
                                                                       z.emplace\_back(c.c + j, p);
                                                                      z.emplace\_back(c.c + k, p);
// remove second
level if to get points for line (defalut: segment) vector<P> CircleCrossLine(P a, P b, P o, double r) {
                                                                    return z;
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y); double C = sq(a.x - o.x)
                                                                  8.12 Area of Union of Circles [0590f1]
                                                                  vector<pair<double , double >>> CoverSegment(C &a , C &b) {
  double d = (a.c - b.c).abs();
       ) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
                                                                    vector<pair<double, double>>> res;
  if (d \ge -eps)
    d = \max(0., d);
                                                                    if (same(a.r + b.r, d))
                                                                    else if (d \le abs(a.r - b.r) + eps) {
    if (a.r < b.r) res.emplace_back(0, 2 * pi);
} else if (d < abs(a.r + b.r) - eps) {
    double i = (-B - \operatorname{sqrt}(d)) / (2 * A);
    double j = (-B + sqrt(d)) / (2 * A);
if (i - 1.0 \le eps \&\& i >=
                                                                      -eps) t.emplace\_back(a.x + i * x, a.y + i * y);
    if (j - 1.0 \le eps \& j > 
                                                                       if (z < 0) z += 2 * pi;
         -eps) t.emplace_back(a.x + j * x, a.y + j * y);
                                                                      double 1 = z - o, r = z + o;

if (1 < 0) 1 += 2 * pi;

if (r > 2 * pi) r -= 2 * pi;
  return t:
                                                                      if (l > r) res.emplace_back (l, 2 * pi), res.emplace_back(0, r);
// calc area
     intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
                                                                      else res.emplace_back(l, r);
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
                                                                    return res;
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
                                                                  double CircleUnionArea
    (vector <C> c) { // circle should be identical
                                                                    int n = c.size();
  if (inb) return
                                                                    double a = 0, w;
        SectorArea(p[0], a, r) + abs(p[0] ^ b)
                                                                    for (int i = 0; w = 0, i < n; ++i) {
                                                                      vector<pair<double, double>>> s = {{2 * pi, 9}}, z; for (int j = 0; j < n; +j) if (i != j) {
     (p.size() = 2u) return SectorArea(a, p[0], r)
        + SectorArea(p[1], b, r) + abs(p[0]
                                                  p[1]) / 2;
                                                                         z = CoverSegment(c[i], c[j]);
  else return SectorArea(a, b, r);
                                                                         for (auto &e : z) s.push_back(e);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
                                                                      sort(s.begin(), s.end());
                                                                      auto F = [\&] (double t) { return c[i].r * (c[i].r *
  double ans = 0;
```

t + c[i].c.x \* sin(t) - c[i].c.y \* cos(t));};

```
for (auto &e : s) {
  if (e.first > w) a += F(e.first) - F(w);
    w = max(w, e.second);
return a * 0.5;
```

#### 8.13 Minimun Distance of 2 Polygons [e9c988]

```
// p, q is convex double TwoConvexHullMinDist
      (Point P[], Point Q[], int n, int m) {
   int YMinP = 0, YMaxQ = 0;
   for (i =
          0; i < n; ++i) if(P[i].y < P[YMinP].y) YMinP = i;
   for (i =
          0; i < m; ++i) if(Q[i].y > Q[YMaxQ].y) YMaxQ = i;
   P[\, n\,] \; = P[\, 0\,] \; , \; \, Q[m] \; = Q[\, 0\,] \, ;
   for (int i = 0; i < n; ++i) {
      while (tmp = Cross(
             \begin{array}{l} \text{Q[YMaxQ + 1] - P[YMinP + 1], P[YMinP] - P[YMinP + 1])} > \text{Cross}\left(\text{Q[YMaxQ] - P[YMinP + 1], P[YMinP} \right) \end{array} 
      if (tmp < 0) ans = min(ans, PointToSegDist (P[YMinP], P[YMinP + 1]), Q[YMaxQ]));
       \begin{array}{l} \textbf{else} \ \ ans = \min(ans \,, \ TwoSegMinDist(P[ \\ YMinP] \,, \ P[YMinP \,+ \,\, 1] \,, \ Q[YMaxQ] \,, \ Q[YMaxQ \,+ \,\, 1])) \,; \end{array} 
      YMinP = (YMinP + 1) \% n;
   }
   return ans;
```

#### 8.14 2D Convex Hull [d97646]

```
bool operator < (const P &a, const P &b) {
  return same(a.x, b.x) ? a.y < b.y : a.x < b.x;
bool operator > (const P &a, const P &b) {
  #define crx(a, b, c) ((b - a) \hat{} (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  \mathtt{sort}\,(\,\mathtt{ps.begin}\,(\,)\,\,,\,\,\,\mathtt{ps.end}\,(\,)\,\,,\,\,\,[\,\&\,]\,\,\,(\,\mathtt{P}\,\,\mathtt{a}\,,\,\,\,\mathtt{P}\,\,\mathtt{b}\,)\,\,\,\{\,\,\,\mathtt{return}\,\,
        same(a.x, b.x) ? a.y < b.y : a.x < b.x; });
    or (int i = 0; i < ps.size(); ++i) {
while (p.size() >= 2 \&\& crx(p[p.size() -
         2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() -
          [2], ps[i], p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
 p.\,pop\_back\,(\,)\;;
  return p;
int sgn(double
      x) \{ return same(x, 0) ? 0 : x > 0 ? 1 : -1; \}
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector < P > p, u, d;
  CH() {}
 CH(vector < P > ps) : p(ps) {
    n = ps.size();
    rotate (p. begin
         (), min_element(p.begin(), p.end()), p.end());
    auto t = max_element(p.begin(), p.end());
d = vector<P>(p.begin(), next(t));
    u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find (vector <P> &v, P d) {
    int l = 0, r = v.size();
     while (1 + 5 < r) {
```

```
else l = L;
     int x = 1;
     for (int i = l +
            1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
     return x;
   int findFarest(P v) {
     if (v.y > 0 | v.y = 0 \& v.x > 0) return
            ((int)d.size() - 1 + find(u, v)) % p.size();
     return find(d, v);
  Pget(int 1, int r, Pa, Pb) {
     int s = sgn(crx(a, b, p[1 \% n]));
     while (l + 1 < r) {
        int m = (l + r) >> 1;
        if (\operatorname{sgn}(\operatorname{crx}(a, b, p[m \% n])) == s) l = m;
        else r = m;
     return isLL(a, b, p[1 \% n], p[(l + 1) \% n]);
   vector <P> getLineIntersect (Pa, Pb) {
     int X = \text{findFarest}((b - a).\text{rot}(pi / 2));

int Y = \text{findFarest}((a - b).\text{rot}(pi / 2));

if (X > Y).\text{swap}(X, Y);
     if (sgn
            \begin{array}{l} (\operatorname{crx}(a,\ b,\ p[X])) \ * \ \operatorname{sgn}(\operatorname{crx}(a,\ b,\ p[Y])) < 0) \\ \operatorname{return} \ \{ \operatorname{get}(X,\ Y,\ a,\ b) \,,\ \operatorname{get}(Y,\ X+n,\ a,\ b) \}; \end{array} 
     return {}; // tangent case falls here
   void update_tangent(P q, int i, int &a, int &b) {
     \begin{array}{ll} if & (sgn(crx(q,\ p[a]\,,\ p[i])) > 0)\ a = i\,;\\ if & (sgn(crx(q,\ p[b],\ p[i])) < 0)\ b = i\,; \end{array}
   void bs(int 1, int r, Pq, int &a, int &b) {
     if (l = r) return;
     update_tangent(q, 1 % n, a, b);
     int s = sgn(crx(q, p[1 \% n], p[(1 + 1) \% n]));

while (1 + 1 < r) {
   int m = (1 + r) >> 1;
        if (sgn(crx
             (q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
        else r = m;
     update_tangent(q, r % n, a, b);
   int x = 1;
  for (int i = l)
         + 1; i < r; ++i) if (v[i] * d > v[x] * d) x = i;
  return x;
int findFarest(P v) {
   if (v.y > 0 | v.y = 0 \& v.x > 0) return
         ((int)d.size() - 1 + find(u, v)) \% p.size();
  return find(d, v);
  get(int 1, int r, Pa, Pb) {
  int s = sgn(crx(a, b, p[1 \% n]));
   while (l + 1 < r) {
     int m = (l + r) >> 1;
      \mbox{if } (sgn(crx(a,\ b,\ p[m\ \%\ n])) == s) \ l = m; \\
     else r = m;
  return isLL(a, b, p[1 % n], p[(1 + 1) % n]);
vector <P> getIS (Pa, Pb) {
   \begin{array}{ll} \text{int } X = findFarest((b-a).spin(pi / 2));\\ \text{int } Y = findFarest((a-b).spin(pi / 2)); \end{array} 
   if (X > Y) swap(X, Y)
   if (\operatorname{sgn}(\operatorname{crx}(a, b, p[X])) * \operatorname{sgn}(\operatorname{crx}(a, b, p[Y])) <
        0) return \{ get(X, Y, a, b), get(Y, X + n, a, b) \};
  return { };
void update_tangent(P q, int i, int &a, int &b) { if (sgn(crx(q, p[a], p[i])) > 0) a = i;
  if (sgn(crx(q, p[b], p[i])) < 0) b = i;
void bs(int l, int r, Pq, int &a, int &b) {
  if (l = r) return;
  update_tangent(q, 1 % n, a, b);
    int \ s = sgn(crx(q, \ p[l \ \% \ n], \ p[(l + 1) \ \% \ n])); 
   while (1 + 1 < r)
     int m = (1 + r) >> 1;
```

```
if (sgn
            (crx(q, p[m \% n], p[(m + 1) \% n])) == s) l = m;
      \begin{array}{ll} \textbf{else} & r = m; \end{array}
   update_tangent(q, r % n, a, b);
bool contain (P p) {
   if (p.x < d[0].x | | p.x > d.back().x) return 0;
   auto it
        = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
   if (it->x = p.x) {
   \begin{array}{l} \mbox{if } (\mbox{it-}>\mbox{y}>\mbox{p.y}) \mbox{ return } 0; \\ \mbox{else if } (\mbox{crx}(\mbox{*prev}(\mbox{it}),\mbox{*it},\mbox{p}) < \mbox{-eps}) \mbox{ return } 0; \end{array}
   it = lower_bound
         (u.begin(), u.end(), P(p.x, 1e12), greater<P>());
   if (it->x = p.x) {
   if (it->y < p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;
bool get_tangent(P p, int &a, int &b) { // b -> a
  if (contain(p)) return 0;
   a = b = 0;
   int i
        = lower_bound(d.begin(), d.end(), p) - d.begin();
  \begin{array}{l} bs(\,0\,,\ i\,,\ p\,,\ a\,,\ b\,)\,;\\ bs(\,i\,,\ d\,.\,size\,(\,)\,,\ p\,,\ a\,,\ b\,)\,; \end{array}
   i = lower\_bound(
         u.begin(), u.end(), p, greater<P>()) - u.begin();
   bs((int
         d.size() - 1, (int)d.size() - 1 + i, p, a, b);
   bs((int)d.size()
            1 + i, (int)d.size() - 1 + u.size(), p, a, b);
};
```

#### 8.15 3D Convex Hull [clae8f]

```
absvol(const P a, const P b, const P c, const P d) {
  return abs(((b-a)^(c-a))*(d-a))/6;
struct convex3D {
  static const int maxn=1010;
  struct T{
    int a,b,c;
    bool res;
    T()\{\}
    T(int a, int
          b, int c, bool res=1: a(a), b(b), c(c), res(res){}
  int n,m;
 P p [maxn];
 T f [maxn*8];
  int id [maxn] [maxn];
  bool on (T &t, P &q) {
    return ((
         p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
  int g=id[a][b];
    if (f[g].res){
       if(on(f[g],p[q]))dfs(q,g);
         id [q][b]=id[a][q]=id[b][a]=m;
         f[m++]=T(b,a,q,1);
    }
  void dfs(int p,int i){
    f[i].res=0;
    meow(p, f[i].b, f[i].a);

meow(p, f[i].c, f[i].b);
    meow(p, f[i].a, f[i].c);
  void operator()(){
    if (n<4)return;
    if ([&](){
         for (int i=1;i< n;++i) if (abs
             (p[0]-p[i])>eps)return swap(p[1],p[i]),0;
         return 1;
}() || [&](){
         for (int
                  i=2; i < n; ++i) if (abs((p[0]-p[i])
              ^(p[1]-p[i])>eps)return swap(p[2],p[i]),0;
         return 1
         }() || [&](){
```

```
for (int i
                                                       = 3; i < n; ++i) if (abs(((p[1]-p[0])^(p[2]-p[0]))
                                                       *(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
                                     }())return;
                   for (int i=0; i<4;++i){
                          T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
                            if(on(t,p[i]))swap(t.b,t.c);
                            id [t.a][t.b]=id [t.b][t.c]=id [t.c][t.a]=m;
                            f[m++]=t;
                   for (int i=4; i< n; ++i) for
                                     (int j=0; j < m++j) if (f[j].res & on(f[j],p[i])) {
                            dfs(i,j);
                            break;
                   int mm=m; m=0;
                   for (int i=0; i \triangleleft mm + +i) if (f[i].res) f[m++]=f[i];
          bool same(int i, int j)
                  return !(absvol(p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].c]),p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].b])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a])>eps || absvol(p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[i].a],p[f[
                                     (p[f[i].a], p[f[i].b], p[f[i].c], p[f[j].c]) > eps);
          int faces(){
                   int r=0;
                   for (int i=0; i \triangleleft m; ++i){
                            int iden=1;
                            for (int j=0; j< i; ++j) if (same(i,j)) iden=0;
                           r += iden;
                   return r;
} tb;
```

### 8.16 Minimum Enclosing Circle [7e5b31]

```
pt center (const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^
                    p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {
        (norm2(cent - p[i]) \le r) continue;
     cent = p[i];
     r = 0.0;
     for (int j = 0; j < i; ++j) {
       if (norm2(cent - p[j]) \le r) continue;
       cent = (p[i] + p[j]) / r = norm2(p[j] - cent);
       for (int k = 0; k < j; ++k) {
         if (norm2(cent - p[k]) <= r) continue;</pre>
         cent = center(p[i], p[j], p[k]);
         r = norm2(p[k] - cent);
    }
  }
  return circle(cent, sqrt(r));
```

#### 8.17 Closest Pair [7f292a]

```
double closest_pair(int l, int r) {
  // p should be sorted
        increasingly according to the x-coordinates.
     (1 = r) return 1e9;
  if (r - l = 1) return dist(p[l], p[r]);
  int m = (l + r) >> 1;
  double d =
       min(closest_pair(l, m), closest_pair(m + 1, r));
  vector<int> vec;
  \quad \text{for (int } i = m; \ i >= 1 \ \&\& \\
       fabs(p[m].x - p[i].x) < d; --i) \ vec.push\_back(i);
      (int i = m + 1; i \le r \&\&
       fabs(p[m].x - p[i].x) < d; ++i) vec.push_back(i);
  sort(vec.begin(), vec.end()
        [\&](int \ a, \ int \ b) \ \{ \ return \ p[a].y < p[b].y; \ \});
  for (int i = 0; i < vec.size(); ++i) {
```

```
for (int j = i + 1; j < vec.size()
    && fabs(p[vec[j]].y - p[vec[i]].y) < d; ++j) {
    d = min(d, dist(p[vec[i]], p[vec[j]]));
}
}
return d;
}</pre>
```

#### 9 Else

#### 9.1 Cyclic Ternary Search\* [28a883]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

# 9.2 Mo's Algorithm (With modification) $_{[5dec12]}$

```
Mo's Algorithm With modification
 Block: N^{2/3}, Complexity: N^{5/3}
 struct Query
    int L, R, LBid, RBid, T;
Query(int l, int r, int t):
    L(1), R(r), LBid(1 / blk), RBid(r / blk), T(t) {}
bool operator < (const Query &q) const {
        \quad \text{if } (LBid \mathrel{!=} q.LBid) \enspace \textbf{return} \enspace LBid < q.LBid;
        if (RBid != q.RBid) return RBid < q.RBid;
       return T < b.T;
 };
 void solve(vector<Query> query) {
    sort(ALL(query));
    int L=0, R=0, T=-1;
    \begin{array}{lll} & for \ (auto \ q \ : \ query) \ \{ \\ & while \ (T < q.T) \ addTime(L, \ R, \ +\!\!+\!\!T) \, ; \ /\!/ \ TODO \end{array}
        while (T > q.T) subTime(L, R, T--); // TODO
       while (R < q.R) add(arr[++R]); // TODO while (L > q.R) add(arr[--L]); // TODO while (R > q.R) sub(arr[R--]); // TODO while (L < q.L) sub(arr[L++]); // TODO
        // answer query
}
```

#### 9.3 Mo's Algorithm On Tree [4a7f74]

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord [in[u]] = ord [out[u]] = u
4) bitset MAXN> inset
struct Query
  int L, R, LBid, lca;
  Query(int u, int v) {
     int c = LCA(u, v);
     if (c = u \mid \mid c = v)
       q.\,lc\,a\,=\,-1\,,\,\,q.\,L\,=\,out\,[\,c\,\,\widehat{}\,\,u\,\,\widehat{}\,\,v\,]\,\,,\,\,q.\,R\,=\,out\,[\,c\,\,]\,;
     else if (out[u] < in[v])
       q.lca = c, q.L = out[u], q.R = in[v];
     else
       q.lca = c, q.L = out[v], q.R = in[u];
    q.Lid = q.L / blk;
  bool operator < (const Query &q) const {
    if (LBid != q.LBid) return LBid < q.LBid;
    return R < q.R;
void flip(int x) {
    if (inset[x]) sub(arr[x]); // TODO
     else add(arr[x]); // TODO
```

```
inset [x] = ~inset [x];
}
void solve(vector<Query> query) {
    sort(ALL(query));
    int L = 0, R = 0;
    for (auto q : query) {
        while (R < q.R) flip(ord[++R]);
        while (L > q.L) flip(ord[--L]);
        while (R > q.R) flip(ord[R--]);
        while (L < q.L) flip(ord[L++]);
        if (~q.lca) add(arr[q.lca]);
        // answer query
        if (~q.lca) sub(arr[q.lca]);
    }
}</pre>
```

#### 9.4 Additional Mo's Algorithm Trick

- Mo's Algorithm With Addition Only
  - Sort queryssame as the normal Mo's algorithm.
  - For each query [l,r]:
  - If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk + 1) \cdot blk$ , curR := curL 1
  - If r > curR, increase curR
  - decrease curL to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
  - Part1: Answer all f([1,r],r+1) first.
  - Part2: Store  $curR \rightarrow R$  for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
  - Note: You must do the above symmetrically for the left boundaries.

#### 9.5 Hilbert Curve [ed5979]

#### 9.6 DynamicConvexTrick\* [6a6f6d]

```
only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
      <(const Line &rhs) const { return a < rhs.a; }</pre>
  bool operator <(ll x) const { return p < x; }
struct DynamicHull : multiset<Line, less >> {
  static const ll kInf = 1e18;
  if (y == end()) { x->p = kInf; return 0; }
    if (x
        ->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
    else x->p = Div(y->b - x->b, x->a - y->a);
    return x->p>=y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
    () && isect(--x, y)) isect(x, y = erase(y));
while ((y = x) != begin
        () && (-x)-p >= y-p isect(x, erase(y));
  ll query(ll x) {
    auto l = *lower_bound(x);
    return l.a * x + l.b;
};
```

#### 9.7 All LCS\* [ae68f0]

### 9.8 AdaptiveSimpson\* [dc2085]

```
template<typename Func, typename d = double>
struct Simpson {
  \begin{array}{ll} \textbf{using} \hspace{0.1cm} \textbf{pdd} \hspace{0.1cm} = \hspace{0.1cm} \textbf{pair} \hspace{-0.1cm} < \hspace{-0.1cm} \textbf{d} \hspace{0.1cm} ; \end{array}
  Func f;
  d eval(pdd l, pdd r, d fm, d eps) {
    pdd m((1.X + r.X) / 2, fm);
     d = mix(1, r, fm) . second;
    auto [flm, sl] = mix(l, m);
      auto [fmr, sr] = mix(m, r); 
     d \ delta = sl + sr - s;
     if (abs(delta
          ) <= 15 * eps) return sl + sr + delta / 15;
     return eval(1, m, flm, eps / 2) +
       eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
     return eval
          (\{l, f(l)\}, \{r, f(r)\}, f((l+r) / 2), eps);
  d eval2(d l, d r, d eps, int k = 997) {
  d h = (r - l) / k, s = 0;
     for (int i = 0; i < k; ++i, l + s; s + eval(1, 1 + h, eps / k);
                                        l += h
     return s;
};
template<typename Func>
Simpson<Func> make_simpson(Func f) { return {f}; }
```

#### 9.9 Simulated Annealing [b14262]

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
    if (exp(-(nw - ans
        ) / factor) >= (double)(rnd() % base) / base)
        ans = nw;
    factor *= 0.99995;
}
```

# 9.10 Tree Hash\* [e57357]

```
ull seed;
ull shift(ull x) {
    x ^= x << 13;
    x ^= x >> 7;
    x ^= x << 17;
    return x;
}
ull dfs(int u, int f) {
    ull sum = seed;
    for (int i : G[u])
        if (i != f)
        sum += shift(dfs(i, u));
    return sum;
}</pre>
```

#### 9.11 Binary Search On Fraction [951597]

```
struct Q {
    ll p, q;
    Q go(Q b, ll d) { return {p + b.p*d, q + b.q*d}; }
};
bool pred(Q);
```

```
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
   Q lo{0, 1}, hi{1, 0};
   if (pred(lo)) return lo;
   assert(pred(hi));
   bool dir = 1, L = 1, H = 1;
   for (; L || H; dir = !dir) {
      ll len = 0, step = 1;
      for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)
        if (Q mid = hi.go(lo, len + step);
            mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
      else len += step;
      swap(lo, hi = hi.go(lo, len));
      (dir ? L : H) = !!len;
   }
   return dir ? hi : lo;
}
```

# 9.12 Bitset LCS [a82d86]

```
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)
   cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {
   cin >> x, (g = f) |= p[x];
   f.shiftLeftByOne(), f.set(0);
   ((f = g - f) ^= g) &= g;
}
cout << f.count() << '\n';</pre>
```

# 10 Python

#### 10.1 Misc