

The choice of which algorithm to use can be critical

ITERATION RECURSION

Recursion

- < Form sorting algorithms
- < Set up programs using pseudocode
- < Use algorithm analysis to see which routine is the best

The great trick in recursion is
that a function is calling itself

pip

"pip installs Python"
or "pip installs packages"

5
4
3
2
1
0

```
1 def countdown(n):  
2     print (n)  
3     if n > 0:  
4         countdown(n-1)  
5
```

To Print Numbers from n Down to 0

```
< Print number n  
< Print numbers from  
n-1 down to 0
```

Divide and conquer

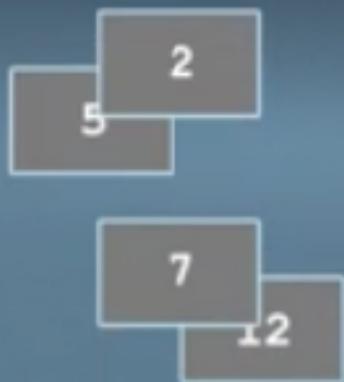
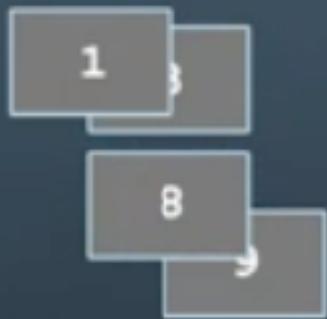
Taking a large data set and
dividing it into subsets that
we handle independently

Merge sort

3
9
8
1

5
2
12
7

Merge sort



Merge sort



```
#Input: unsorted list L, length n  
#Output: L, with elements sorted smallest to largest  
  
If n <= 1  
    Return L  
    #a list of length 1 is already sorted
```

Base case

```
#Input: unsorted list L, length n
#Output: L, with elements sorted smallest to largest

If n <= 1
    Return L
    #a list of length 1 is already sorted
L1 = L(0:n/2-1) L2 = (n/2:n-1)
MergeSort(L1) MergeSort(L2)
L = Merge(L1, L2)
#Merge will be defined separately
```

```
1 def mergeSort(L):
2     n = len(L)
3     if n <= 1:
4         return
5     L1 = L[:n//2]
6     L2 = L[n//2:]
7     mergeSort(L1)
8     mergeSort(L2)
9     merge(L, L1, L2)
10    return
11
12    If n <= 1
13        Return L
14
15        L1 = L(0:n/2-1) L2 = (n/2:n-1)
16        MergeSort(L1) MergeSort(L2)
17        L = Merge(L1, L2)
```

```
1 def merge(L, L1, L2):
2     i = 0
3     j = 0
4     k = 0
5     while (j < len(L1)) or (k < len(L2)):
6         if j < len(L1):
7             if k < len(L2):
8                 # we are not at the end of L1 or L2, so pull the smaller value
9                 if L1[j] < L2[k]:
10                     L[i] = L1[j]
11                     j += 1
12                 else:
13                     L[i] = L2[k]
14                     k += 1
15             else:
16                 # we are at the end of L2, so just pull from L1
17                 L[i] = L1[j]
18                 j += 1
19         else:
20             # we are at the end of L1, so just pull from L2
21             L[i] = L2[k]
22             k += 1
23         i += 1
24     return
```



```
['barbeque', 'chicken and dumplings', 'gumbo', 'ice cream', 'pecan pie', 'pizza']
```

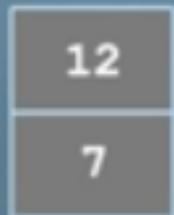
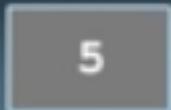
```
1 favorite_foods = ['pizza', 'barbeque', 'gumbo', 'chicken and dumplings', 'pecan pie', 'ice cream']
2 mergeSort(favorite_foods)
3 print(favorite_foods)
```

Quick sort

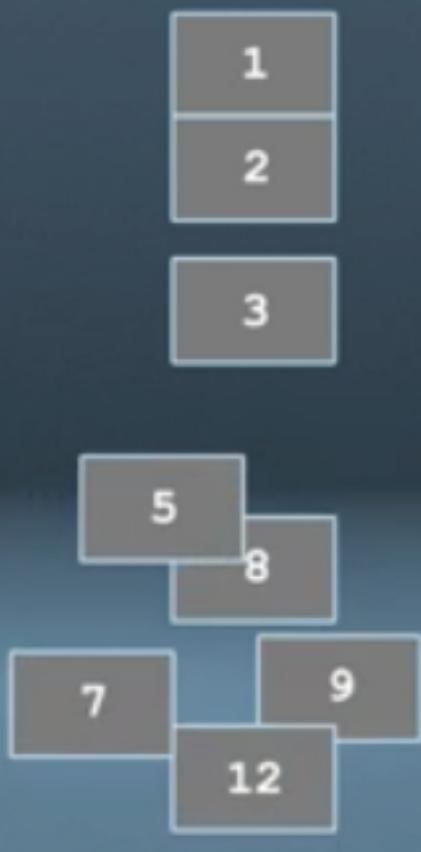
9
8
1
5
2
12
7

3 Pivot

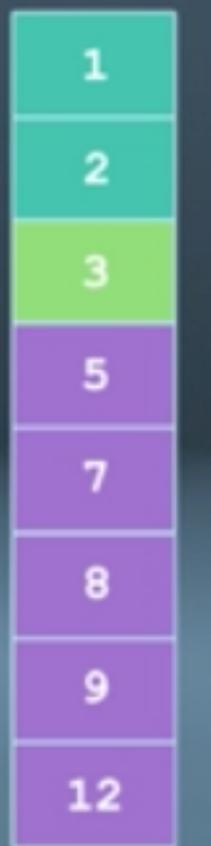
Quick sort



Quick sort



Quick sort



Quick Sort

```
< Input: unsorted list L, length n  
< Output: L, with elements sorted smallest  
to largest
```



Quick Sort

```
< Input: unsorted list L, length n  
< Output: L, with elements sorted smallest  
to largest  
1. If n <= 1  
   1. Return L  
      #a list of length 1 is already sorted
```



Quick Sort

```
< Input: unsorted list L, length n
< Output: L, with elements sorted smallest
  to largest
1. If n <= 1
  1. Return L
    #a list of length 1 is already sorted
2. pivot = L[0]
3. Use remaining elements to form list L1 less
  than, and L2 greater than, the pivot
```



Quick Sort

```
< Input: unsorted list L, length n
< Output: L, with elements sorted smallest
  to largest
1. If n <= 1
  1. Return L
    #a list of length 1 is already sorted
2. pivot = L[0]
3. Use remaining elements to form list L1 less
  than, and L2 greater than, the pivot
4. QuickSort(L1) QuickSort(L2)
```



Quick Sort

```
< Input: unsorted list L, length n
< Output: L, with elements sorted smallest
  to largest
1. If n <= 1
  1. Return L
    #a list of length 1 is already sorted
2. pivot = L[0]
3. Use remaining elements to form list L1 less
  than, and L2 greater than, the pivot
4. QuickSort(L1) QuickSort(L2)
5. L = Join(L1, pivot, L2)
```



```
1 def quickSort(L):
2     #handle base case
3     if len(L) <= 1:
4         return
5
6     #pick pivot
7     pivot = L[0]
8
9     #form lists less/greater than pivot
10    L1 = []
11    L2 = []
12    for element in L[1:]:
13        if element < pivot:
14            L1.append(element)
15        else:
16            L2.append(element)
17
18    #sort sublists
19    quickSort(L1)
20    quickSort(L2)
21
22    #join the sublists and pivot
23    L[:] = []
24    for element in L1:
25        L.append(element)
26    L.append(pivot)
27    for element in L2:
28        L.append(element)
29
30    return
```

```
'barbeque', 'chicken and dumplings', 'gumbo', 'ice cream', 'pecan pie', 'pizza'
```

```
1 favorite_foods = ['pizza', 'barbeque', 'gumbo', 'chicken and dumplings', 'pecan pie', 'ice cream']
2 quickSort(favorite_foods)
3 print(favorite_foods)
4
```

Sort Routines

- < Selection sort
- < Insertion sort
- < Merge sort
- < Quick sort

Choosing Between Algorithms

- < Easiest for programmer to understand
- < Easiest to write code for
- < Increased efficiency

Asymptotic analysis

Looks at how a function performs as the input size grows larger and larger

"What will happen if I double the input size?"

Linear Search

n items

- < Best case; first item in the list
- < Worst case; go through all n items and its not in the list, or its last
- < Average case; half the items ($n/2$)

Linear Search

$2n$ items

- < Best case; first item in the list
- < Worst case; go through all $2n$ items and its not in the list
- < Average case; half the items ($2n/2=n$)

Linear Search $2n$ items

- < Best case; first item in the list
 - < Worst case; go through all $2n$ items and its not in the list
 - < Average case; half the items ($2n/2=n$)
- Linear search; $O(n)$

Binary Search

- < Best case: first item in the list
- < Worst case: keep dividing by 2 until down to a single element

Size of List	Max Number of Iterations
1	1
2	2
3	2
4	3
5	3
6	3
7	3
8	4
9-15	4
16-31	5
32-63	6
64-127	7
128-255	8
256-511	9
512-1023	10
2^{n-1} through $2^n - 1$	n

Log function

Running time is a logarithmic function of the input size

Binary Search

- < Best case: first item in the list

- < Worst case: keep dividing by 2 until down to a single element

$O(\log_2 n)$

If choosing between a linear process and a logarithmic process, given a sufficiently large data set, always prefer the log process

selection sort

63
18
42
7
99
26
71

selection sort

63
18
42
7
99
26
71

Selection Sort

< n

< n-1

< n-2

< n-3

< ...

< 1

< Summation formula: $\frac{n(n+1)}{2}$

< Quadratic running time: $O(n^2)$

Quick Sort

< Input: unsorted list L, length n

< Output: L, with elements sorted smallest to largest

1. If $n \leq 1$
 1. Return L
#a list of length 1 is already sorted
2. pivot = L[0]
3. Use remaining elements to form list L1 less than, and L2 greater than, the pivot
4. QuickSort(L1) QuickSort(L2)
5. L = Join(L1, pivot, L2)

Insertion sort



Best Case

$O(n)$

Worst Case

1

2

3

...

n

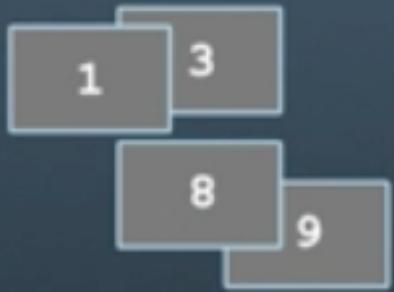
$O(n^2)$

Merge sort

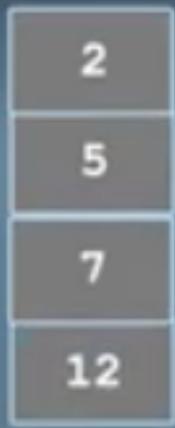
3
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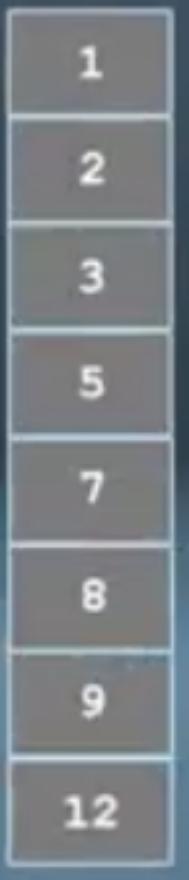
Merge sort



Merge sort



Merge sort



Operation	Best	Average	Worst
Search	$O(1)$	$O(n)$	$O(n)$
Binary Search	$O(1)$	$O(\log n)$	$O(\log n)$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quicksort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$

sort() Function

- < Sorting algorithm in Python
- < Designed to be more efficient than selection sort, insertion sort, merge sort and quicksort
- < Combination of merge sort and insertion sort

sort() Function

- < Uses merge sort for the overall problem
- < Switches to insertion sort
- < Checks to see if portions of the list are already sorted

Fibonacci Numbers

$$< F(0) = 0, \quad F(1) = 1$$

$$< F(n) = F(n-2) + F(n-1)$$

$$< 0$$

$$1$$

$$(0+1) = 1$$

$$(1+1) = 2$$

$$(1+2) = 3$$

$$(2+3) = 5$$

$$(3+5) = 8$$

...

```
1 def Fib(n):
2     if n == 0:
3         return 0
4     elif n == 1:
5         return 1
6     else:
7         return Fib(n-2) + Fib(n-1)
8
```

```
1 Fib(100) =  
2 Fib(98) + Fib(99) =  
3 (Fib(96) + Fib(97)) + (Fib(97) + Fib(98)) =  
4 ((Fib(94) + Fib(95)) + (Fib(95)+ Fib(96))) +  
5 ((Fib(95) + Fib(96)) + (Fib(96) + Fib(97)))  
6 ...  
7  
8  
9      2^100 function calls  
10
```

Nonillion
trillion * trillion * trillion

exponential
time

10 million
years

```
1 Fib(100) =  
2 Fib(98) + Fib(99) =  
3 (Fib(96) + Fib(97)) + (Fib(97) + Fib(98)) =  
4 ((Fib(94) + Fib(95)) + (Fib(95)+ Fib(96))) +  
5 ((Fib(95) + Fib(96)) + (Fib(96) + Fib(97)))  
6 ...  
7  
8  
9  
10  
11  
12  
13  
14  
15
```



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```
1 def Fib2(n):
2     F = [0,1]
3     for i in range(2,n+1):
4         #compute F[i]
5         F.append(F[-2] + F[-1])
6     return F[n]
7 print(Fib2(100))
8
```

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linear algorithm

```
1 def Fib2(n):
2     F = [0,1]
3     for i in range(2,n+1):
4         #compute F[i]
5         F.append(F[-2] + F[-1])
6     return F[n]
7 print(Fib2(100))
8
```

It's worth analyzing your code to determine running time, to make sure you're not being unnecessarily inefficient in the algorithm you've chosen to solve a problem