CSPs of Finite Commutative Idempotent Binars

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joint work with

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Shanks Workshop Vanderbilt University May 30, 2015

slides available at

https://github.com/williamdemeo/Talks

CSP DICHOTOMY CONJECTURE

For a (finite, idempotent) algebra A...

CSP(A) is tractable \iff A has a weak-nu term operation

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For a (finite, idempotent) algebra A...

CSP(A) is tractable \implies A has a weak-nu term operation \checkmark

The left-to-right direction is known.

CSP DICHOTOMY CONJECTURE

For a (finite, idempotent) algebra A...

CSP(A) is tractable \iff A has a weak-nu term operation (?)

The right-to-left direction is open.

CSP DICHOTOMY CONJECTURE

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A weak near unanimity (weak-nu) term operation is one that satisfies

$$t(x, x, \dots, x) \approx x$$
 (idempotent)

$$t(y, x, \dots, x) \approx t(x, y, \dots, x) \approx \dots \approx t(x, x, \dots, y)$$

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 $t(y, x, ..., x) \approx t(x, y, ..., x) \approx ... \approx t(x, x, ..., y)$

A binary operation t(x, y) is weak-nu if

$$t(x,x) \approx x$$
 (idempotent) $t(y,x) \approx t(x,y)$ (commutative)

So let's try to prove (?) for commutative idempotent binars.

A CIB is an algebra $\mathbf{A} = \langle A, \cdot \rangle$ satisfying $x \cdot y \approx y \cdot x$ and $x \cdot x \approx x$.

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First Example: a semilattice is an associative CIB.

Semilattices are tractable.

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First Example: a semilattice is an associative CIB. Semilattices are tractable.

Pause to consider more general case for a minute...

SOME WELL KNOWN FACTS

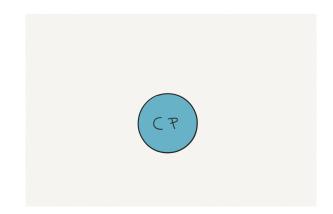
Let A be a finite idempotent algebra. Let S_2 be the 2-elt semilattice.

V(A) is CP \iff A has Malcev term

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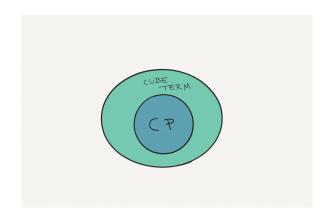
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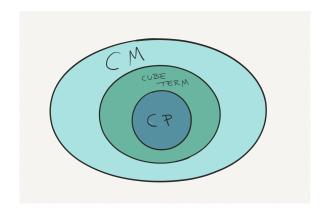
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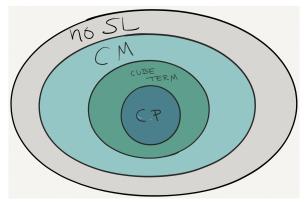
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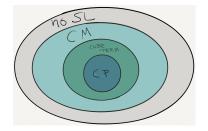
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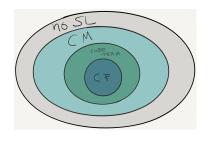
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 $\mathbf{A} = \mathbf{a}$ finite idempotent algebra $\mathbf{S}_2 = \mathbf{the}$ 2-elt semilattice.

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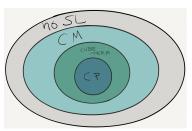
- (Berman, Idziak, Marković, McKenzie, Valeriote, Willard 2010)
- CM \Longrightarrow S₂ is not in V(A) Proof: S₂ ∈ V(A) \Rightarrow S₂² ∈ V(A); Con (S₂²) is not modular.

■ cube term ⇒ CM

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CIB case

1st reduction by cube-term blockers.

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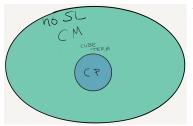
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Marković, M. Maróti, McKenzie (M^4) "Finitely related clones and algebras with cube terms" (2012)

A cube-term blocker (CTB) is a pair (C,B) of subuniverses satisfying $\emptyset < C < B \leqslant A$ and for every $t(x_1,\ldots,x_n)$ there is an index $i \in [n]$ with

$$(\forall (b_1,\ldots,b_n)\in B^n)(b_i\in C\longrightarrow t(b_1,\ldots,b_n)\in C).$$

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LEMMA

A finite CIB $\mathbf A$ has a CTB if and only if $\mathbf S_2 \in \mathsf{HS}(\mathbf A)$.

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LEMMA

A finite CIB A has a CTB if and only if $S_2 \in \mathsf{HS}(A)$.

PROOF.

(C,B) a CTB implies $\theta = C^2 \cup (B-C)^2$ a congruence with $\mathbf{B}/\theta \cong \mathbf{S}_2$.

Conversely, suppose $S_2 \in \mathsf{HS}(\mathbf{A})$, and \mathbf{B} is a subalgebra of \mathbf{A} with \mathbf{B}/θ a meet-SL for some θ . Let C/θ be the bottom of \mathbf{B}/θ , then (C,B) is a CTB.

SECOND REDUCTION

Kearnes and Tschantz

"Automorphism groups of squares and of free algebras" (2007)

LEMMA

If V is an idempotent variety that is not congruence permutable, then there are subuniverses U and W of $\mathbf{F} := \mathbf{F}_V\{x,y\}$ satisfying

- 1. $x \in U \cap W$
- 2. $y \in U^c \cap W^c$
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For CIB's, either *U* or *W* will be an ideal.

This implies a CTB and a semilattice.

CONCLUSION

Let A be a finite CIB. Then

 $\textbf{S}_2\notin \text{HS}(\textbf{A})$ if and only if $\,V(\textbf{A})$ is congruence permutable.

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Recall, if $V(\mathbf{A})$ is SD_{\wedge} , then $CSP(\mathbf{A})$ is tractable.

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Recall, if V(A) is SD_{\wedge} , then CSP(A) is tractable.

REVISED QUESTION

Let A be a finite CIB with S_2 in HS(A), and V(A) not SD_{\wedge} .

Is CSP(A) tractable?

	0	1	2	3
0	0	0	0	1
1	0	1	3	2
2 3	0	3	2	1
3	1	2	1	3

Cliff's idea: replace basic binary operation with a term from $\operatorname{Clo}(\mathbf{A})$, say $t(x,y)=(x\cdot(x\cdot y))\cdot(y\cdot(x\cdot y))$. If $\langle A,t\rangle$ tractable, then so is $\mathbf{A}=\langle A,\cdot\rangle$.

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If
$$\langle A, t \rangle$$
 tractable, then so is $\mathbf{A} = \langle A, \cdot \rangle$.

$$\{t\} \subseteq \operatorname{Clo}(\mathbf{A}) \implies \operatorname{Rel}(\operatorname{Clo}(\mathbf{A})) \subseteq \operatorname{Rel}(\{t\})$$

$$\implies$$
 CSP(A) \leq_P CSP $\langle A, t \rangle$

	0	1	2	3	
0	0	0	0	1	
1	0	1	3	2	
1 2 3	0	3	2	1	
3	1	2	1	3	
		-	_	0	
t	0	1	2	3	
0	0	0	0	0	
0		0	0 3	0	
	0	-	0	0	

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$$\implies \operatorname{CSP}(\mathbf{A}) \leq_{P} \operatorname{CSP}(A, t)$$

 $\langle A, t \rangle$ tractable \implies **A** tractable

	0	1	2	3
0	0	0	1	1
1	0	1	3	2
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Let
$$t_2(x, y) = x \cdot (x \cdot (x \cdot y)) \cdot y \cdot (y \cdot (x \cdot y)).$$

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	0	1	2	3	
0	0	0	1	1	
0 1 2 3	0 0 1 1	1 3 2	3 2	2	
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to	l 0	1	2	3	

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$$\langle A, t_2 \rangle$$
 tractable

•	U	- 1	2	3
0	0	0	2	1
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1 2 3	0 0 2	3 2	2	1
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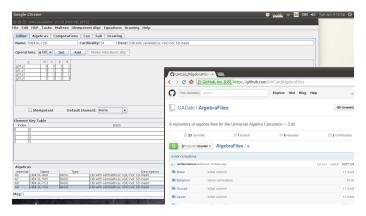
		0	1	2	3
,	0	0	0	2	1
	1	0 0 2	1	3	2
	2	2	3	2	1
	2	4	2	4	2

Let
$$t_3(x, y) = ...$$
 ?

	0	1	2	3
0	0	0	2	1
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0	4	0	4	0

Let
$$t_3(x, y) = ...$$
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Let $t_3(x, y, z) = ...$?

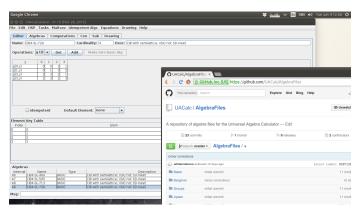
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Thank you for listening!