

DEDEKIND'S TRANSPOSITION PRINCIPLE  
AND  
ISOTOPIC ALGEBRAS WITH NONISOMORPHIC  
CONGRUENCE LATTICES

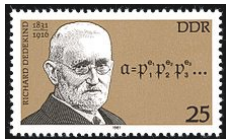
William DeMeo  
williamdemeo@gmail.com

University of South Carolina

AMS Spring Western Sectional Meeting  
University of Colorado, Boulder, CO

April 13-14, 2013

*These slides and other resources are available at*  
<http://williamdemeo.wordpress.com>



# DEDEKIND'S TRANSPOSITION PRINCIPLE

## FOR MODULAR LATTICES

### Notation

Let  $\mathbf{L} = \langle L, \wedge, \vee \rangle$  be a lattice with  $a \in L$ .

Let  $\varphi_a$  and  $\psi_a$  be the *perspectivity maps*

$$\varphi_a(x) = x \wedge a \quad \text{and} \quad \psi_a(x) = x \vee a$$

For  $x, y \in L$ , let  $\llbracket x, y \rrbracket_L = \{z \in L \mid x \leq z \leq y\}$ .

# DEDEKIND'S TRANSPOSITION PRINCIPLE

## FOR MODULAR LATTICES

### Notation

Let  $\mathbf{L} = \langle L, \wedge, \vee \rangle$  be a lattice with  $a \in L$ .

Let  $\varphi_a$  and  $\psi_a$  be the *perspectivity maps*

$$\varphi_a(x) = x \wedge a \quad \text{and} \quad \psi_a(x) = x \vee a$$

For  $x, y \in L$ , let  $\llbracket x, y \rrbracket_L = \{z \in L \mid x \leq z \leq y\}$ .

### THEOREM (DEDEKIND'S TRANSPOSITION PRINCIPLE)

$\mathbf{L}$  is modular iff for all  $a, b \in L$  the maps  $\varphi_a$  and  $\psi_b$  are inverse lattice isomorphisms of  $\llbracket a \wedge b, a \rrbracket$  and  $\llbracket b, a \vee b \rrbracket$ .

# DEDEKIND'S TRANSPOSITION PRINCIPLE

## FOR MODULAR LATTICES

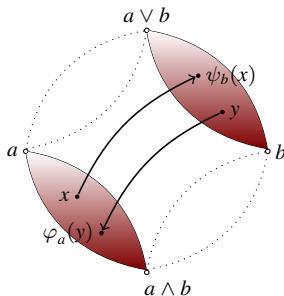
### Notation

Let  $\mathbf{L} = \langle L, \wedge, \vee \rangle$  be a lattice with  $a \in L$ .

Let  $\varphi_a$  and  $\psi_a$  be the *perspectivity maps*

$$\varphi_a(x) = x \wedge a \quad \text{and} \quad \psi_a(x) = x \vee a$$

For  $x, y \in L$ , let  $\llbracket x, y \rrbracket_L = \{z \in L \mid x \leq z \leq y\}$ .



### THEOREM (DEDEKIND'S TRANSPOSITION PRINCIPLE)

$\mathbf{L}$  is modular iff for all  $a, b \in L$  the maps  $\varphi_a$  and  $\psi_b$  are inverse lattice isomorphisms of  $\llbracket a \wedge b, a \rrbracket$  and  $\llbracket b, a \vee b \rrbracket$ .

## ANOTHER TRANSPOSITION PRINCIPLE

### FOR LATTICES OF EQUIVALENCE RELATIONS

Let  $X$  be a set and let  $\text{Eq } X$  be the lattice of equivalence relations on  $X$ .

If  $L$  is a sublattice of  $\text{Eq } X$  with  $\eta, \theta \in L$ , then we define

$$[\![\eta, \theta]\!]_L = \{\gamma \in L \mid \eta \leq \gamma \leq \theta\}.$$

## ANOTHER TRANSPOSITION PRINCIPLE

### FOR LATTICES OF EQUIVALENCE RELATIONS

Let  $X$  be a set and let  $\text{Eq } X$  be the lattice of equivalence relations on  $X$ .

If  $L$  is a sublattice of  $\text{Eq } X$  with  $\eta, \theta \in L$ , then we define

$$[\![\eta, \theta]\!]_L = \{\gamma \in L \mid \eta \leq \gamma \leq \theta\}.$$

For  $\beta \in \text{Eq } X$ , let  $[\![\eta, \theta]\!]_L^\beta$  be the set of relations in  $[\![\eta, \theta]\!]_L$  that permute with  $\beta$ ,

$$[\![\eta, \theta]\!]_L^\beta = \{\gamma \in L \mid \eta \leq \gamma \leq \theta \text{ and } \gamma \circ \beta = \beta \circ \gamma\}.$$

## ANOTHER TRANSPOSITION PRINCIPLE

### FOR LATTICES OF EQUIVALENCE RELATIONS

Let  $X$  be a set and let  $\text{Eq } X$  be the lattice of equivalence relations on  $X$ .

If  $L$  is a sublattice of  $\text{Eq } X$  with  $\eta, \theta \in L$ , then we define

$$[\eta, \theta]_L = \{\gamma \in L \mid \eta \leq \gamma \leq \theta\}.$$

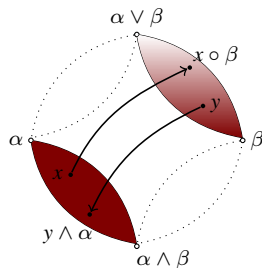
For  $\beta \in \text{Eq } X$ , let  $[\eta, \theta]_L^\beta$  be the set of relations in  $[\eta, \theta]_L$  that permute with  $\beta$ ,

$$[\eta, \theta]_L^\beta = \{\gamma \in L \mid \eta \leq \gamma \leq \theta \text{ and } \gamma \circ \beta = \beta \circ \gamma\}.$$

### LEMMA

Suppose  $\alpha$  and  $\beta$  are permuting relations in  $L \leq \text{Eq } X$ .

Then  $[\beta, \alpha \vee \beta]_L \cong [\alpha \wedge \beta, \alpha]_L^\beta \leq [\alpha \wedge \beta, \alpha]_L$ .



## DEDEKIND'S RULE

The proof requires the following version of *Dedekind's Rule*:

### LEMMA

*Suppose  $\alpha, \beta, \gamma \in L \leq \text{Eq } X$  and  $\alpha \leq \beta$ .*

*Then the following identities of subsets of  $X^2$  hold:*

$$\alpha \circ (\beta \cap \gamma) = \beta \cap (\alpha \circ \gamma)$$

$$(\beta \cap \gamma) \circ \alpha = \beta \cap (\gamma \circ \alpha)$$



# ISOTOPY

## BASIC DEFINITIONS

Let  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  be algebras of the same type.

$\mathbf{A}$  and  $\mathbf{B}$  are *isotopic over*  $\mathbf{C}$ , denoted  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ , if there is an isomorphism

$\varphi : \mathbf{A} \times \mathbf{C} \xrightarrow{\cong} \mathbf{B} \times \mathbf{C}$  that leaves the second coordinate fixed

i.e.  $(\forall a \in A) (\forall c \in C) \quad \varphi(a, c) = (\varphi_1(a, c), c)$

# ISOTOPY

## BASIC DEFINITIONS

Let  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  be algebras of the same type.

$\mathbf{A}$  and  $\mathbf{B}$  are *isotopic over*  $\mathbf{C}$ , denoted  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ , if there is an isomorphism

$\varphi : \mathbf{A} \times \mathbf{C} \xrightarrow{\cong} \mathbf{B} \times \mathbf{C}$  that leaves the second coordinate fixed

$$\text{i.e. } (\forall a \in A) (\forall c \in C) \quad \varphi(a, c) = (\varphi_1(a, c), c)$$

We say that  $\mathbf{A}$  and  $\mathbf{B}$  are *isotopic*, denoted  $\mathbf{A} \sim \mathbf{B}$ , if  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  for some  $\mathbf{C}$ .

It is easy to verify that  $\sim$  is an equivalence relation.

# ISOTOPY

## BASIC DEFINITIONS

Let  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  be algebras of the same type.

$\mathbf{A}$  and  $\mathbf{B}$  are *isotopic over*  $\mathbf{C}$ , denoted  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ , if there is an isomorphism

$\varphi : \mathbf{A} \times \mathbf{C} \xrightarrow{\cong} \mathbf{B} \times \mathbf{C}$  that leaves the second coordinate fixed

$$\text{i.e. } (\forall a \in A) (\forall c \in C) \quad \varphi(a, c) = (\varphi_1(a, c), c)$$

We say that  $\mathbf{A}$  and  $\mathbf{B}$  are *isotopic*, denoted  $\mathbf{A} \sim \mathbf{B}$ , if  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  for some  $\mathbf{C}$ .

If  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  and  $\text{Con}(\mathbf{A} \times \mathbf{C})$  happens to be modular, then we write  $\mathbf{A} \sim_{\mathbf{C}}^{\text{mod}} \mathbf{B}$  and say that  $\mathbf{A}$  and  $\mathbf{B}$  are *modular isotopic over*  $\mathbf{C}$ .

# ISOTOPY

## BASIC DEFINITIONS

Let  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  be algebras of the same type.

$\mathbf{A}$  and  $\mathbf{B}$  are *isotopic over*  $\mathbf{C}$ , denoted  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ , if there is an isomorphism

$\varphi : \mathbf{A} \times \mathbf{C} \xrightarrow{\cong} \mathbf{B} \times \mathbf{C}$  that leaves the second coordinate fixed

$$\text{i.e. } (\forall a \in A) (\forall c \in C) \quad \varphi(a, c) = (\varphi_1(a, c), c)$$

We say that  $\mathbf{A}$  and  $\mathbf{B}$  are *isotopic*, denoted  $\mathbf{A} \sim \mathbf{B}$ , if  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  for some  $\mathbf{C}$ .

If  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  and  $\text{Con}(\mathbf{A} \times \mathbf{C})$  happens to be modular, then we write  $\mathbf{A} \sim_{\mathbf{C}}^{\text{mod}} \mathbf{B}$  and say that  $\mathbf{A}$  and  $\mathbf{B}$  are *modular isotopic over*  $\mathbf{C}$ .

We call  $\mathbf{A}$  and  $\mathbf{B}$  *modular isotopic in one step*, denoted  $\mathbf{A} \sim_1^{\text{mod}} \mathbf{B}$ , if they are modular isotopic over some  $\mathbf{C}$ .

# ISOTOPY

## BASIC DEFINITIONS

Let  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  be algebras of the same type.

$\mathbf{A}$  and  $\mathbf{B}$  are *isotopic over*  $\mathbf{C}$ , denoted  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ , if there is an isomorphism

$\varphi : \mathbf{A} \times \mathbf{C} \xrightarrow{\cong} \mathbf{B} \times \mathbf{C}$  that leaves the second coordinate fixed

$$\text{i.e. } (\forall a \in A) (\forall c \in C) \quad \varphi(a, c) = (\varphi_1(a, c), c)$$

We say that  $\mathbf{A}$  and  $\mathbf{B}$  are *isotopic*, denoted  $\mathbf{A} \sim \mathbf{B}$ , if  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  for some  $\mathbf{C}$ .

If  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$  and  $\text{Con}(\mathbf{A} \times \mathbf{C})$  happens to be modular, then we write  $\mathbf{A} \sim_{\mathbf{C}}^{\text{mod}} \mathbf{B}$  and say that  $\mathbf{A}$  and  $\mathbf{B}$  are *modular isotopic over*  $\mathbf{C}$ .

We call  $\mathbf{A}$  and  $\mathbf{B}$  *modular isotopic in one step*, denoted  $\mathbf{A} \sim_1^{\text{mod}} \mathbf{B}$ , if they are modular isotopic over some  $\mathbf{C}$ .

We call  $\mathbf{A}$  and  $\mathbf{B}$  are *modular isotopic*, denoted  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$ , if  $(\mathbf{A}, \mathbf{B})$  is in the transitive closure of  $\sim_1^{\text{mod}}$ .

# ISOTOPY

## MODULAR CASE

**Lemma 11.** If  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$  then  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ .

The proof is a nice/easy application of Dedekind's Transposition Principle.

# ISOTOPY

## MODULAR CASE

**Lemma 11.** If  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$  then  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ .

The proof is a nice/easy application of Dedekind's Transposition Principle.

Could we use the same strategy with the non-modular version of the transposition principle to show that  $\mathbf{A} \sim \mathbf{B}$  implies  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ ?

# ISOTOPY

## MODULAR CASE

**Lemma 11.** If  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$  then  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ .

The proof is a nice/easy application of Dedekind's Transposition Principle.

Could we use the same strategy with the non-modular version of the transposition principle to show that  $\mathbf{A} \sim \mathbf{B}$  implies  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ ?

As you have guessed, the answer is no!

The perspectivity map that is so useful when  $\text{Con}(\mathbf{A} \times \mathbf{C})$  is modular can fail *miserably* in the non-modular case...



# ISOTOPY

## MODULAR CASE

**Lemma 11.** If  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$  then  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ .

The proof is a nice/easy application of Dedekind's Transposition Principle.

Could we use the same strategy with the non-modular version of the transposition principle to show that  $\mathbf{A} \sim \mathbf{B}$  implies  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ ?

As you have guessed, the answer is no!

The perspectivity map that is so useful when  $\text{Con}(\mathbf{A} \times \mathbf{C})$  is modular can fail *miserably* in the non-modular case... *even when*  $\mathbf{A} \cong \mathbf{B}$ !

# ISOTOPY

## MODULAR CASE

**Lemma 11.** If  $\mathbf{A} \sim^{\text{mod}} \mathbf{B}$  then  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ .

The proof is a nice/easy application of Dedekind's Transposition Principle.

Could we use the same strategy with the non-modular version of the transposition principle to show that  $\mathbf{A} \sim \mathbf{B}$  implies  $\text{Con } \mathbf{A} \cong \text{Con } \mathbf{B}$ ?

As you have guessed, the answer is no!

The perspectivity map that is so useful when  $\text{Con}(\mathbf{A} \times \mathbf{C})$  is modular can fail *miserably* in the non-modular case... *even when*  $\mathbf{A} \cong \mathbf{B}$ !

But this only shows that the same argument doesn't work...

## COUNTEREXAMPLES

We describe a class of examples in which  $\mathbf{A} \sim \mathbf{B}$  and  $\text{Con } \mathbf{A} \not\cong \text{Con } \mathbf{B}$ .

The examples show that congruence lattices of isotopic algebras can differ arbitrarily in size.

## COUNTEREXAMPLES

We describe a class of examples in which  $\mathbf{A} \sim \mathbf{B}$  and  $\text{Con } \mathbf{A} \not\cong \text{Con } \mathbf{B}$ .

The examples show that congruence lattices of isotopic algebras can differ arbitrarily in size.

For any group  $G$ , let  $\text{Sub}(G)$  denote the lattice of subgroups of  $G$ .

## COUNTEREXAMPLES

We describe a class of examples in which  $\mathbf{A} \sim \mathbf{B}$  and  $\text{Con } \mathbf{A} \not\cong \text{Con } \mathbf{B}$ .

The examples show that congruence lattices of isotopic algebras can differ arbitrarily in size.

For any group  $G$ , let  $\text{Sub}(G)$  denote the lattice of subgroups of  $G$ .

A group  $G$  is called a *Dedekind group* if every subgroup of  $G$  is normal.

## COUNTEREXAMPLES

We describe a class of examples in which  $\mathbf{A} \sim \mathbf{B}$  and  $\text{Con } \mathbf{A} \not\cong \text{Con } \mathbf{B}$ .

The examples show that congruence lattices of isotopic algebras can differ arbitrarily in size.

For any group  $G$ , let  $\text{Sub}(G)$  denote the lattice of subgroups of  $G$ .

A group  $G$  is called a *Dedekind group* if every subgroup of  $G$  is normal.

Let  $S$  be any group and let  $D$  denote the *diagonal subgroup* of  $S \times S$ ,

$$D = \{(x, x) \mid x \in S\}$$

The interval  $[[D, S \times S] \leq \text{Sub}(S \times S)$  is described by the following

### LEMMA

*The filter above the diagonal subgroup of  $S \times S$  is isomorphic to the lattice of normal subgroups of  $S$ .*

## THE EXAMPLE

Let  $S$  be a group, and let  $G = S_1 \times S_2$ , where  $S_1 \cong S_2 \cong S$ .

Let  $D = \{(x_1, x_2) \in G \mid x_1 = x_2\}$ ,  $T_1 = S_1 \times \langle 1 \rangle$ ,  $T_2 = \langle 1 \rangle \times S_2$ .

## THE EXAMPLE

Let  $S$  be a group, and let  $G = S_1 \times S_2$ , where  $S_1 \cong S_2 \cong S$ .

Let  $D = \{(x_1, x_2) \in G \mid x_1 = x_2\}$ ,  $T_1 = S_1 \times \langle 1 \rangle$ ,  $T_2 = \langle 1 \rangle \times S_2$ .

Then  $D \cong T_1 \cong T_2$ , and these are pair-wise compliments:

$$\langle T_1, T_2 \rangle = \langle T_1, D \rangle = \langle D, T_2 \rangle = G$$

$$T_1 \cap D = D \cap T_2 = T_1 \cap T_2 = \langle (1, 1) \rangle$$



## THE EXAMPLE

Let  $S$  be a group, and let  $G = S_1 \times S_2$ , where  $S_1 \cong S_2 \cong S$ .

Let  $D = \{(x_1, x_2) \in G \mid x_1 = x_2\}$ ,  $T_1 = S_1 \times \langle 1 \rangle$ ,  $T_2 = \langle 1 \rangle \times S_2$ .

Then  $D \cong T_1 \cong T_2$ , and these are pair-wise compliments:

$$\langle T_1, T_2 \rangle = \langle T_1, D \rangle = \langle D, T_2 \rangle = G$$

$$T_1 \cap D = D \cap T_2 = T_1 \cap T_2 = \langle (1, 1) \rangle$$

Let  $\mathbf{A} = \langle G/T_1, G^{\mathbf{A}} \rangle$  = the algebra with universe the left cosets of  $T_1$  in  $G$ , and basic operations the left multiplications by elements of  $G$ .

For each  $g \in G$  the operation  $g^{\mathbf{A}} \in G^{\mathbf{A}}$  is defined by

$$g^{\mathbf{A}}(xT_1) = (gx)T_1 \quad (xT_1 \in G/T_1).$$

Define the algebra  $\mathbf{C} = \langle G/T_2, G^{\mathbf{C}} \rangle$  similarly.

## THE EXAMPLE

The algebra  $\mathbf{B}$  will have universe  $B = G/D$ , but we define the action of  $G$  on  $B$  with a twist.

## THE EXAMPLE

The algebra  $\mathbf{B}$  will have universe  $B = G/D$ , but we define the action of  $G$  on  $B$  with a twist.

For each  $g = (g_1, g_2) \in G$ , for each  $(x_1, x_2)D \in G/D$ , define

$$g^{\mathbf{B}}((x_1, x_2)D) = (g_2x_1, g_1x_2)D.$$

Let  $\mathbf{B} = \langle G/D, G^{\mathbf{B}} \rangle$ , where  $G^{\mathbf{B}} = \{g^{\mathbf{B}} \mid g \in G\}$ .

## THE EXAMPLE

The algebra  $\mathbf{B}$  will have universe  $B = G/D$ , but we define the action of  $G$  on  $B$  with a twist.

For each  $g = (g_1, g_2) \in G$ , for each  $(x_1, x_2)D \in G/D$ , define

$$g^{\mathbf{B}}((x_1, x_2)D) = (g_2x_1, g_1x_2)D.$$

Let  $\mathbf{B} = \langle G/D, G^{\mathbf{B}} \rangle$ , where  $G^{\mathbf{B}} = \{g^{\mathbf{B}} \mid g \in G\}$ .

Consider the binary relation  $\varphi \subseteq (A \times C) \times (B \times C)$  that associates to each ordered pair

$$((x_1, x_2)T_1, (y_1, y_2)T_2) \in A \times C$$

the pair

$$((x_2, y_1)D, (y_1, y_2)T_2) \in B \times C$$

## THE EXAMPLE

The algebra  $\mathbf{B}$  will have universe  $B = G/D$ , but we define the action of  $G$  on  $B$  with a twist.

For each  $g = (g_1, g_2) \in G$ , for each  $(x_1, x_2)D \in G/D$ , define

$$g^{\mathbf{B}}((x_1, x_2)D) = (g_2x_1, g_1x_2)D.$$

Let  $\mathbf{B} = \langle G/D, G^{\mathbf{B}} \rangle$ , where  $G^{\mathbf{B}} = \{g^{\mathbf{B}} \mid g \in G\}$ .

Consider the binary relation  $\varphi \subseteq (A \times C) \times (B \times C)$  that associates to each ordered pair

$$((x_1, x_2)T_1, (y_1, y_2)T_2) \in A \times C$$

the pair

$$((x_2, y_1)D, (y_1, y_2)T_2) \in B \times C$$

It is easy to verify that this relation is a function, and in fact

$\varphi: \mathbf{A} \times \mathbf{C} \rightarrow \mathbf{B} \times \mathbf{C}$  is an isomorphism.

## THE EXAMPLE

The algebra  $\mathbf{B}$  will have universe  $B = G/D$ , but we define the action of  $G$  on  $B$  with a twist.

For each  $g = (g_1, g_2) \in G$ , for each  $(x_1, x_2)D \in G/D$ , define

$$g^{\mathbf{B}}((x_1, x_2)D) = (g_2x_1, g_1x_2)D.$$

Let  $\mathbf{B} = \langle G/D, G^{\mathbf{B}} \rangle$ , where  $G^{\mathbf{B}} = \{g^{\mathbf{B}} \mid g \in G\}$ .

Consider the binary relation  $\varphi \subseteq (A \times C) \times (B \times C)$  that associates to each ordered pair

$$((x_1, x_2)T_1, (y_1, y_2)T_2) \in A \times C$$

the pair

$$((x_2, y_1)D, (y_1, y_2)T_2) \in B \times C$$

It is easy to verify that this relation is a function, and in fact

$\varphi: \mathbf{A} \times \mathbf{C} \rightarrow \mathbf{B} \times \mathbf{C}$  is an isomorphism.

Since  $\varphi$  leaves second coordinates fixed,  $\mathbf{A} \sim_{\mathbf{C}} \mathbf{B}$ .

## CONCLUSION

Compare Con **A** and Con **B**.

## CONCLUSION

Compare  $\text{Con } \mathbf{A}$  and  $\text{Con } \mathbf{B}$ .

$\text{Con } \mathbf{A} \cong \llbracket T_1, G \rrbracket \leq \text{Sub}(G)$ , so  $\text{Con } \mathbf{A} \cong \text{Sub}(S)$ .



## CONCLUSION

Compare  $\text{Con } \mathbf{A}$  and  $\text{Con } \mathbf{B}$ .

$\text{Con } \mathbf{A} \cong \llbracket T_1, G \rrbracket \leq \text{Sub}(G)$ , so  $\text{Con } \mathbf{A} \cong \text{Sub}(S)$ .

$\text{Con } \mathbf{B}$  is isomorphic to the lattice of normal subgroups of  $S$ .

## CONCLUSION

Compare  $\text{Con } \mathbf{A}$  and  $\text{Con } \mathbf{B}$ .

$\text{Con } \mathbf{A} \cong \llbracket T_1, G \rrbracket \leq \text{Sub}(G)$ , so  $\text{Con } \mathbf{A} \cong \text{Sub}(S)$ .

$\text{Con } \mathbf{B}$  is isomorphic to the lattice of normal subgroups of  $S$ .

$$\text{Con } \mathbf{B} \cong \text{NSub}(S) \leq \text{Sub}(S) \cong \text{Con } \mathbf{A}$$

So, if  $S$  is any non-Dedekind group,  $\text{Con } \mathbf{B} \not\cong \text{Con } \mathbf{A}$ .

## CONCLUSION

Compare  $\text{Con } \mathbf{A}$  and  $\text{Con } \mathbf{B}$ .

$\text{Con } \mathbf{A} \cong \llbracket T_1, G \rrbracket \leq \text{Sub}(G)$ , so  $\text{Con } \mathbf{A} \cong \text{Sub}(S)$ .

$\text{Con } \mathbf{B}$  is isomorphic to the lattice of normal subgroups of  $S$ .

$$\text{Con } \mathbf{B} \cong \text{NSub}(S) \leq \text{Sub}(S) \cong \text{Con } \mathbf{A}$$

So, if  $S$  is any non-Dedekind group,  $\text{Con } \mathbf{B} \not\cong \text{Con } \mathbf{A}$ .

If  $S$  is a nonabelian simple group, then  $\text{Con } \mathbf{B} \cong \mathbf{2}$ , while  $\text{Con } \mathbf{A} \cong \text{Sub}(S)$  can be arbitrarily large.