

# MATH 654 HOMEWORK #2 (DUE ON 9/16/2016)

Your name:

**Problem 8.** Prove that a set,  $X$ , is finite iff there is a relation such that it and its inverse order both well-order  $X$ .

**Definition 2.** A set of strings over alphabet  $A$  is a subset of  $\bigcup_{n \in \omega} A^n$ . In other words, it is a set of functions whose domains are natural numbers and whose ranges are contained in  $A$ . We denote by  $\lambda$  the unique function with domain 0. The alphabet will not be mentioned if it is clear from context or does not need to be specified. If  $S$  is a set of strings and  $x, y \in S$ , we write  $x \preceq y$  if  $y \upharpoonright \text{dom}(x) = x$ . If  $\neg(x \preceq y \vee y \preceq x)$ , then we write  $x \perp y$  and say that  $x$  and  $y$  are incomparable; otherwise, we write  $x \parallel y$  and say that  $x$  and  $y$  are comparable. For a set of strings,  $S$ ,  $T[S] = \{x : (\exists y \in S)(x \preceq y)\}$  is the prefix closure of  $S$ .

**Definition 3.** Let  $S$  and  $P$  be two sets of strings.

- $P * S = \{xy : x \in P \wedge y \in S\}$ .
- $P^{-1}S = \{y : (\exists x \in P)(xy \in S)\}$ .

For notational simplicity, we define  $x^{-1}S = \{x\}^{-1}S$ ,  $P^{-1}x = P^{-1}\{x\}$ ,  $x * S = \{x\} * S$  and  $P * x = P * \{x\}$  for a string  $x$ .

**Definition 4.** Given a set of strings,  $S$ , we call  $P \subseteq T[S]$  a *maximal antichain* of  $S$  if  $(\forall x, y \in P)(x \perp y \vee x = y)$  and  $(\forall x \in S)(\exists y \in P)(y \parallel x)$ .  $P$  is a *valid antichain* of  $S$  if  $P$  is a maximal antichain of  $S$  and  $(\forall x, y \in P)(x^{-1}T[S] = y^{-1}T[S])$ . We define,  $\text{Vac}(S) = \{P : P \text{ is a valid antichain of } S\}$ .

**Example 1.** Consider the following set of strings over the alphabet  $\{a, b\}$ :

$$S = \{a^5, a^4b, a^2ba, a^2b^2, ba^4, ba^3b, baba, bab^2, b^2a^3, b^2a^2b, b^3a, b^4\}.$$

Graphically, we can represent  $S$  as a tree where branching left indicates an  $a$  and branching right indicates a  $b$ . In the picture below to the right, we highlight the four valid antichains of  $S$ :  $P_0 = \{\lambda\}$ ,  $P_1 = \{a^2, ba, b^2\}$ ,  $P_2 = \{a^4, a^2b, ba^3, bab, b^2a^2, b^3\}$  and  $P_3 = S$ . Note that  $S$  is only a valid antichain of itself because it contains no comparable strings. The members of the four valid antichains are connected via dotted lines in the right picture ( $P_0$  has only one member and therefore includes no dotted lines). For reference a maximal antichain that is not valid is included in the picture on the left and its members are joined with a dotted line.

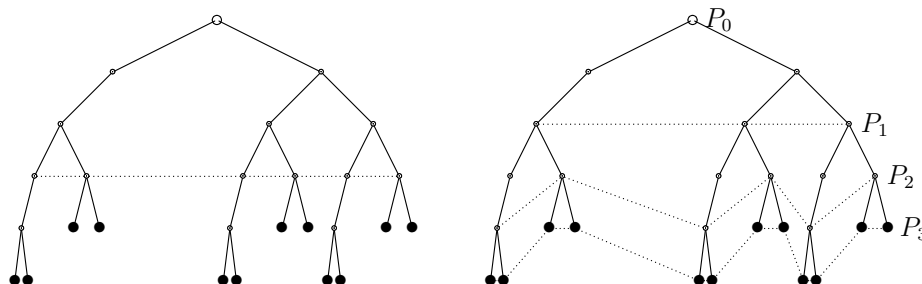


FIGURE 1. On the left, a maximal antichain that is not valid; on the right, all the valid antichains.

In the next figure, we focus on the valid antichain  $P_1$ .

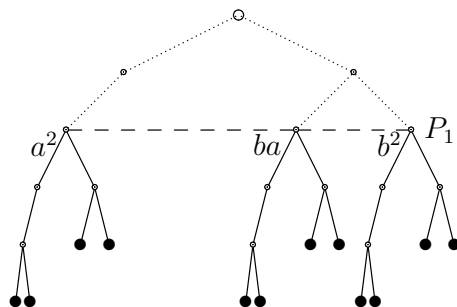


FIGURE 2. The identical subtrees below the elements of the valid antichain  $P_1$ .

Observe that the portions of the tree below each of  $a^2$ ,  $ba$  and  $b^2$  are identical; the terminal nodes of all three sub-trees are  $\{a^3, a^2b, ab, b^2\}$ . It is this equivalence of suffixes that makes  $P_1$  a valid antichain.

**Problem 9.** Suppose that  $P$  is a valid antichain of a set of strings  $S$  and  $Q$  is a valid antichain of  $P$ . Prove that  $Q$  is a valid antichain of  $S$ .