

# **SC2001/CE2101/CZ2101: Algorithm Design and Analysis**

## **Mergesort**

**Instructor: Assoc. Prof. ZHANG Hanwang**

Courtesy of Dr. Ke Yiping, Kelly's slides

# Learning Objectives

At the end of this lecture, students should be able to:

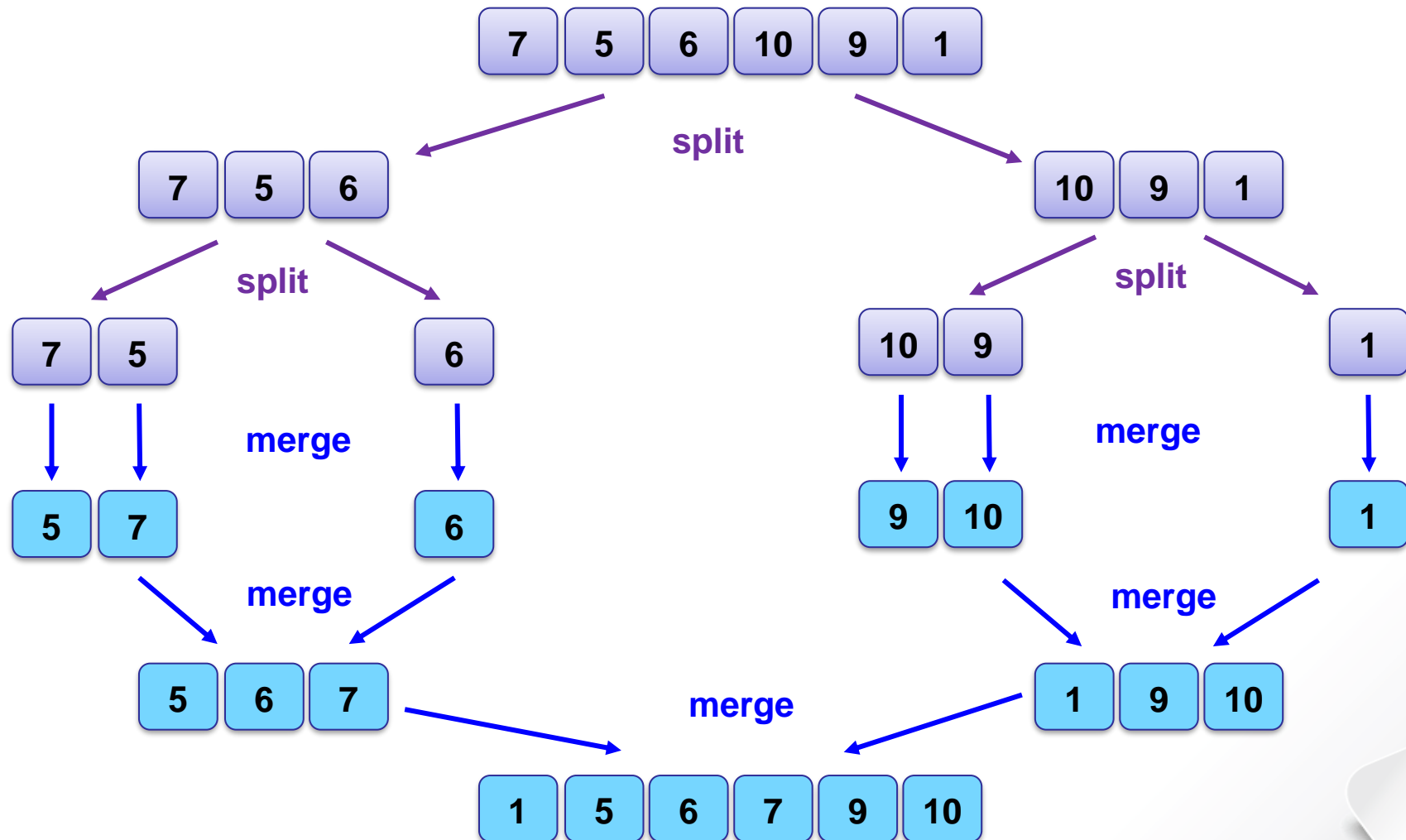
- Explain the approach of **Divide and Conquer**
- Describe how Mergesort works by:
  - Recalling the pseudo code
  - Manually executing the algorithm on a toy input array
- Analyse the **time complexity** of Mergesort, by using:
  - Recurrence equation
  - Recursion tree



# **Mergesort**

## **(Divide and Conquer Approach)**

# Mergesort in a nutshell



# Mergesort

## The Divide and Conquer approach

The skeleton of this approach:

**solve (problem of size  $n$ )**

{ if ( $n \leq$  minimum size)

    solve the problem directly;

else {

    divide the problem into  $p_1, p_2, \dots, p_k$ ;

    for each sub-problem  $p_s$

$\text{solution}_s = \text{solve}(p_s)$ ;

    combine all  $\text{solution}_s$ ;

}

}

# Mergesort

## The Divide and Conquer approach

The skeleton of this approach:

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  { if (n <= minimum size)  
    solve the problem directly;  
  else {  
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       $\text{solution}_s = \text{solve}(p_s)$ ;  
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  }  
}
```

# Mergesort

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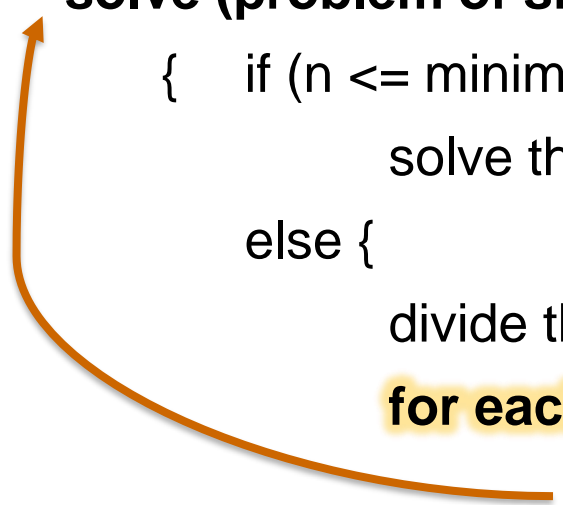
**for each sub-problem  $p_s$**

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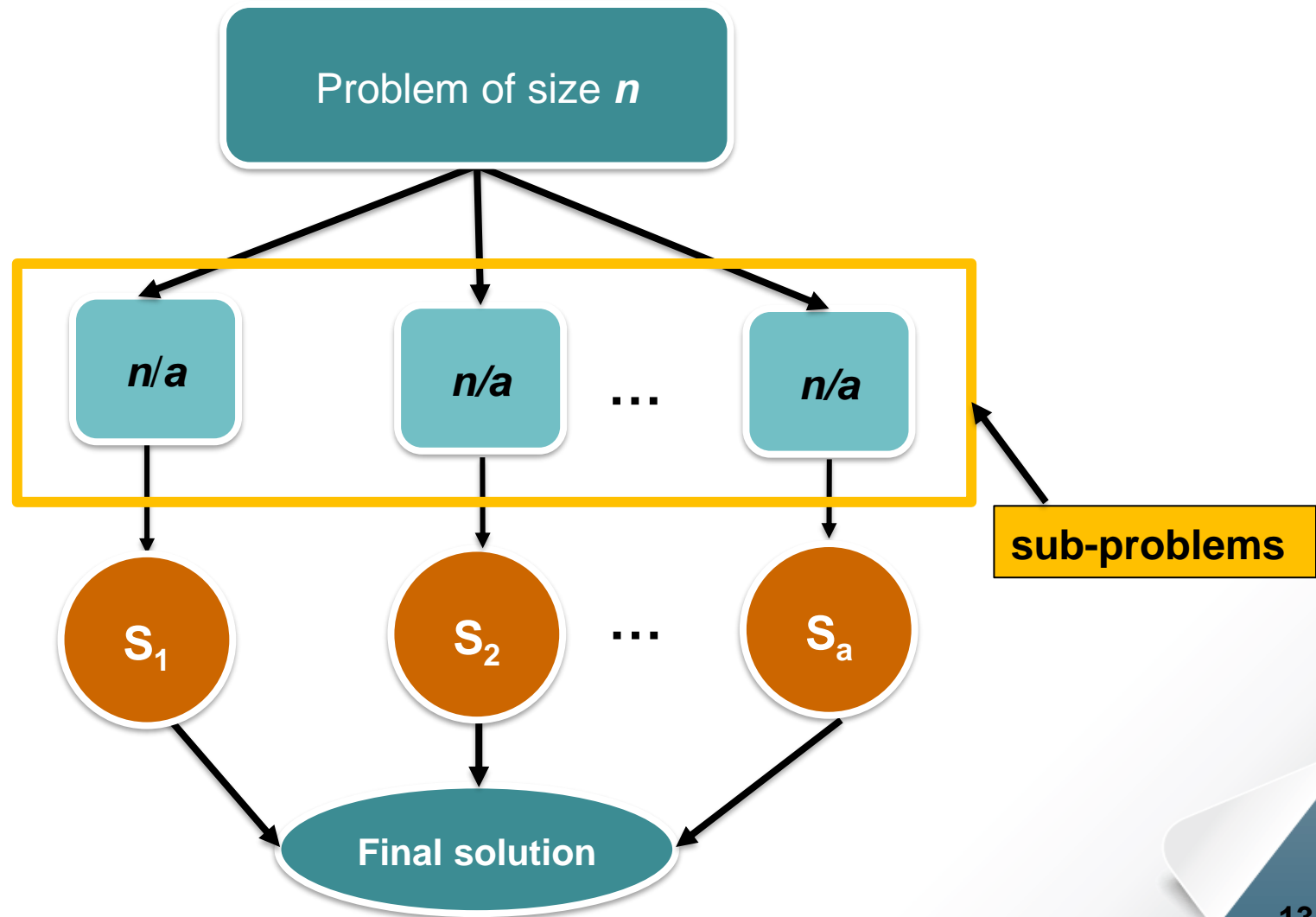
**solution<sub>s</sub> = solve ( $p_s$ );**

**combine all solution<sub>s</sub>;**

    }

}

# Mergesort





# Mergesort (Algorithm)

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```
mergeSort(list) {  
    if (length of list > 1) {  
        Partition list into two (approx.) equal sized  
        lists, L1 & L2;  
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        merge the sorted L1 & L2;  
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```

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        lists, L1 & L2;
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```
        mergeSort (L1);
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```
        mergeSort (L2);
```

```
        merge the sorted L1 & L2;
```

```
    }
```

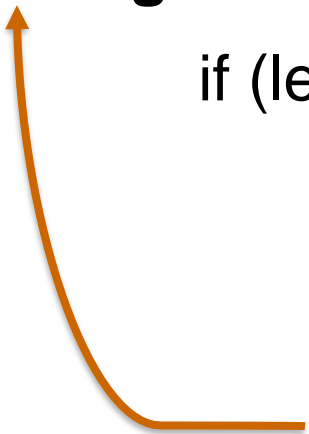
```
}
```

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```

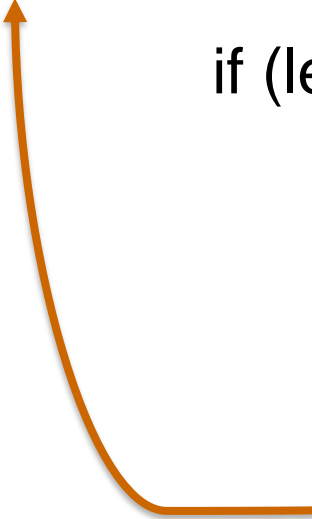
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}
```



# **Mergesort**

## **(Overview of Pseudo Code)**

# Mergesort

```
void mergesort(int n, int m)
```

```
{  int mid = (n+m)/2;
```

```
    if (m-n <= 0)
```

```
        return;
```

```
    else if (m-n > 1) {
```

```
        mergesort(n, mid);
```

```
        mergesort(mid+1, m);
```

```
    }
```

```
    merge(n, m);
```

```
}
```

5	4	3	7	6
---	---	---	---	---





# Mergesort

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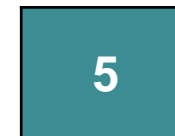
```
  }
```

```
  merge(n, m);
```

```
}
```



if  $m-n = 0$ ,



$m = n$

if  $m-n < 0$ , Empty array



# Mergesort

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void mergesort(int n, int m)
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```
{  int mid = (n+m)/2;
```

```
  if (m-n <= 0)
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    return;
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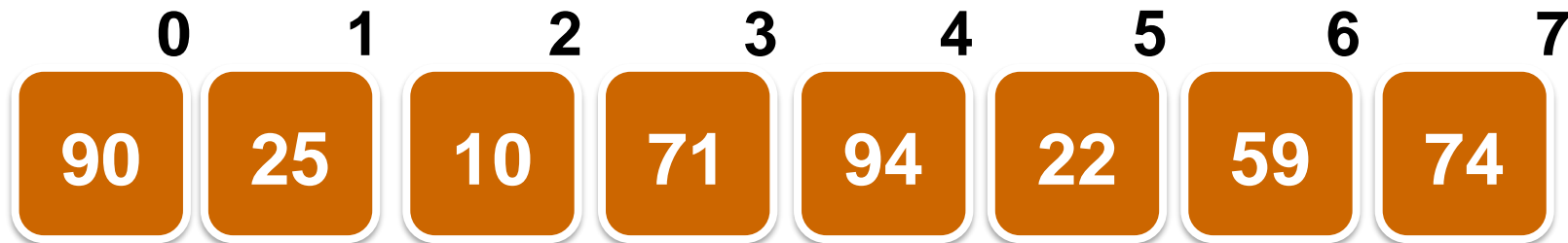




## Mergesort (Example)

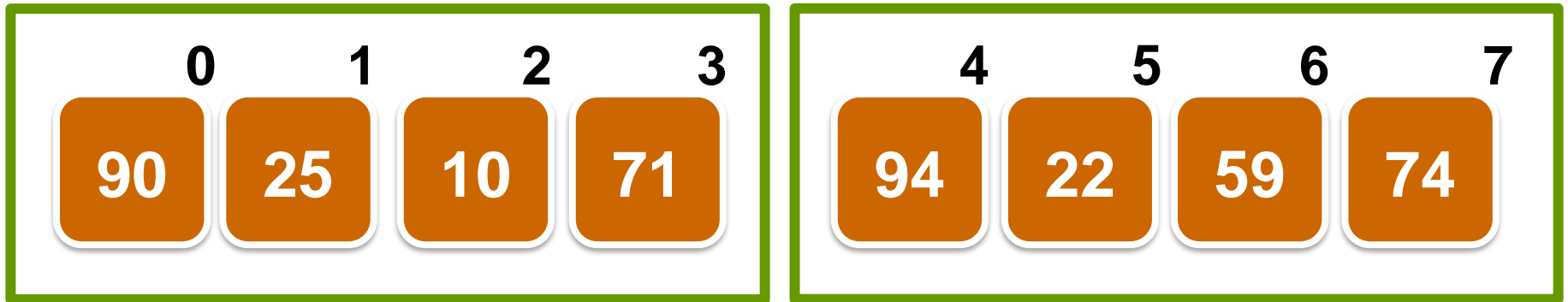
# Mergesort

Sort in ascending order



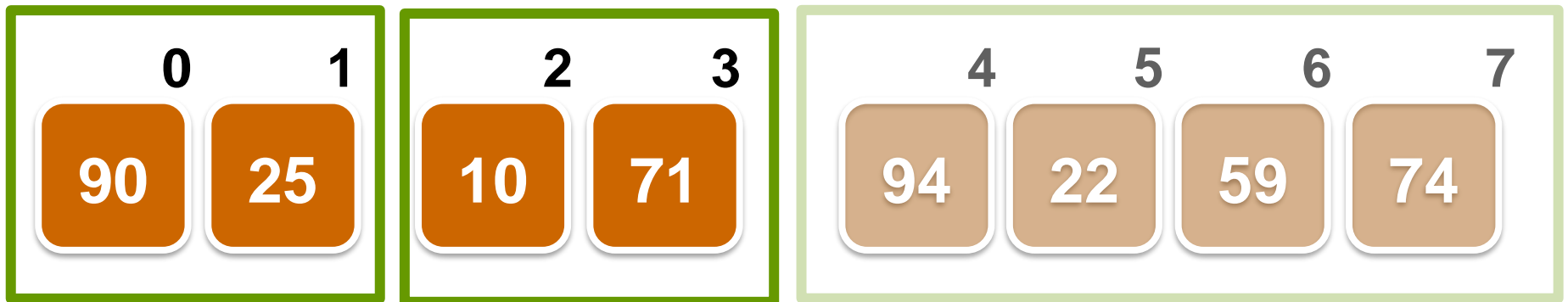
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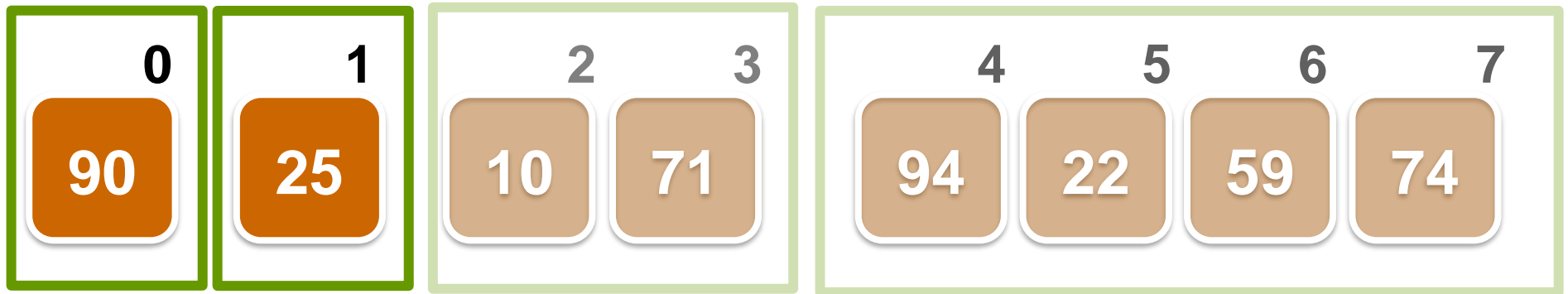
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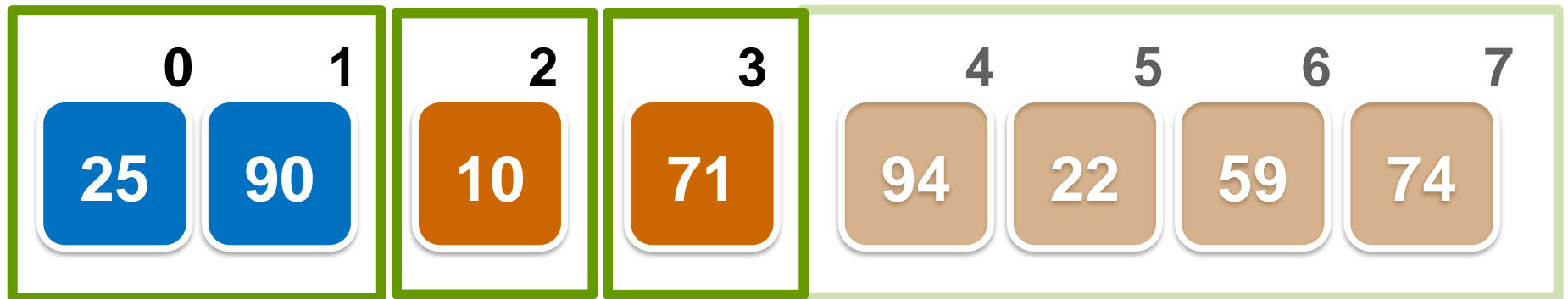
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1 key comparison in merging

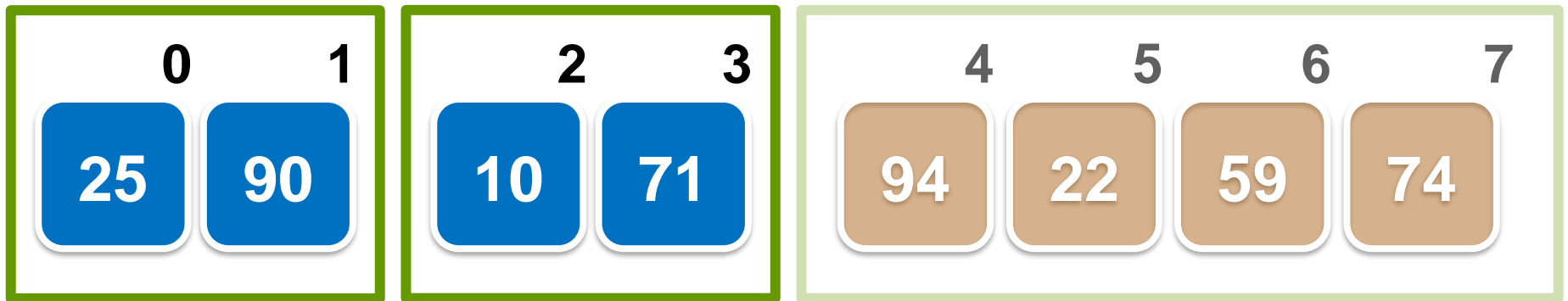
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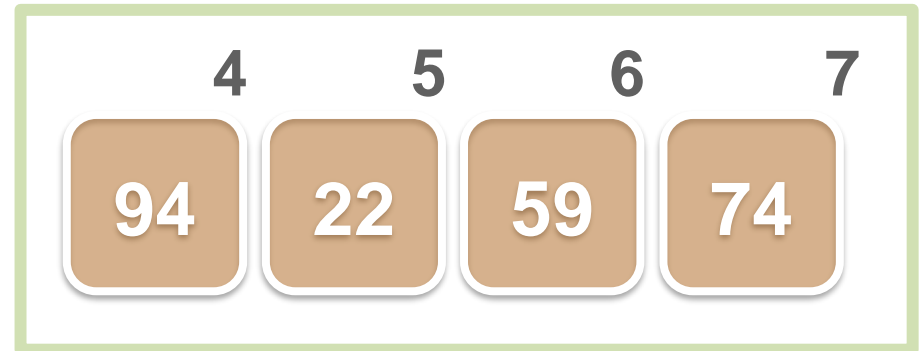
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1 key comparison in merging

# Mergesort

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3 key comparison in merging

# Mergesort

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3 key comparison in merging

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Sort in ascending order





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1 key comparison in merging

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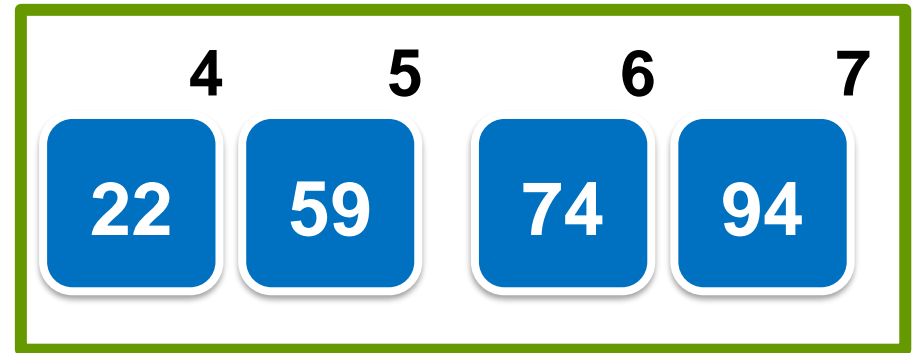
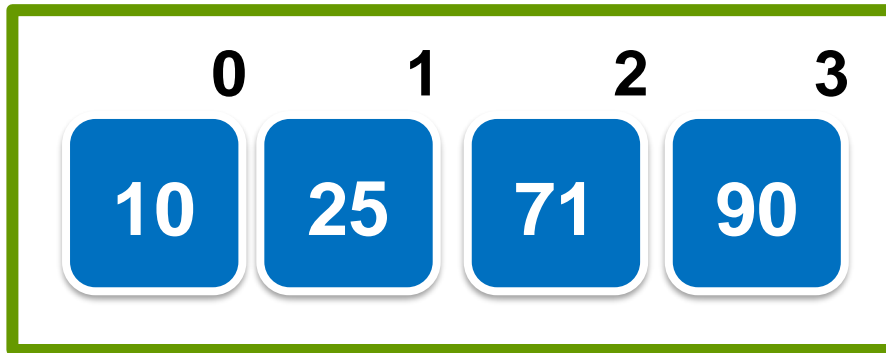
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1 key comparison in merging

# Mergesort

Sort in ascending order

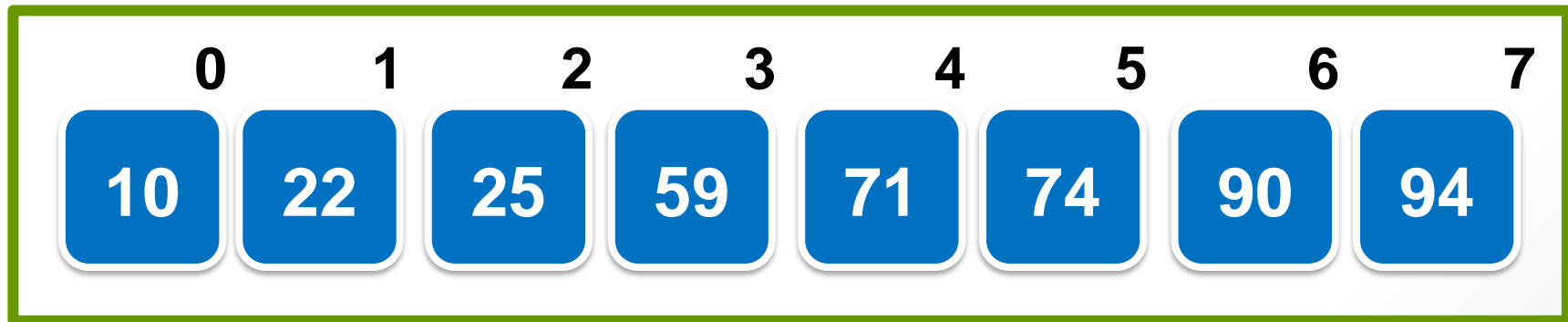


3 key comparison in merging

# Mergesort



Sorted in ascending order



7 key comparisons in merging



# Merge (Pseudo Code)

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void merge(int n, int m) {  
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    while (both halves are not empty) {  
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        if (1st element of 1st half is smaller)  
            1st element of 1st half joins the end of the merged list;  
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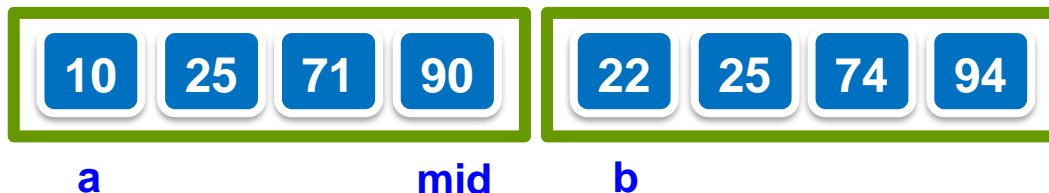
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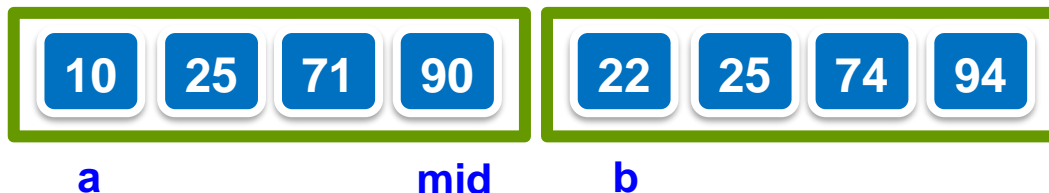
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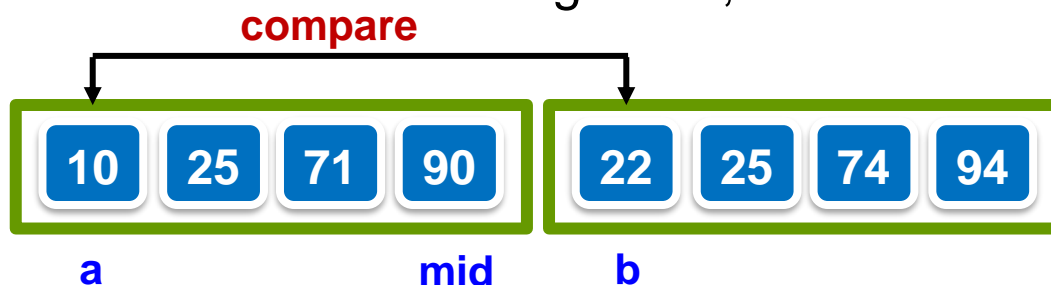
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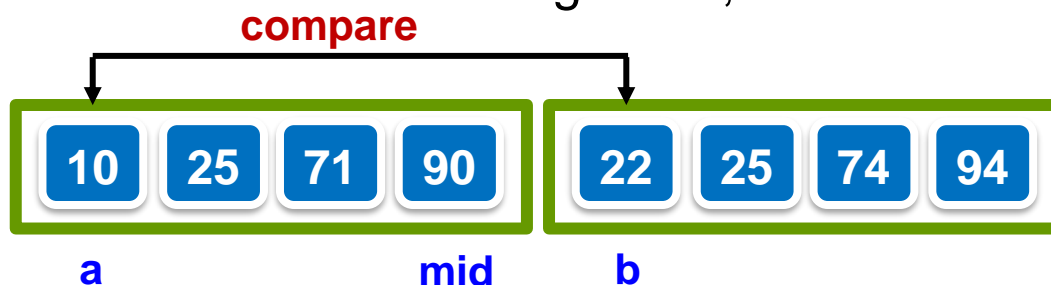
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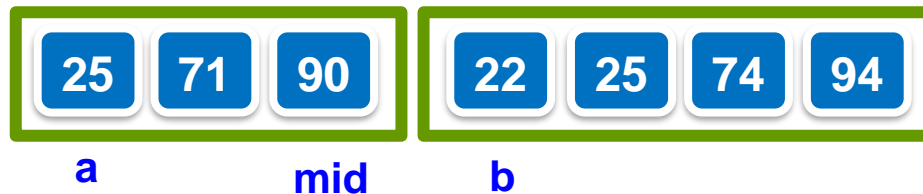
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```

10



# Merge (Pseudo Code)

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}
```



# Merge (Pseudo Code)

```
else { // the 1st elements of the 2 halves are equal
```

```
    if (they are the last elements) break;
```

```
        1st element of 1st half joins end of the merged list;
```

```
        move the 1st element of 2nd half to the end of the  
        merged list;
```

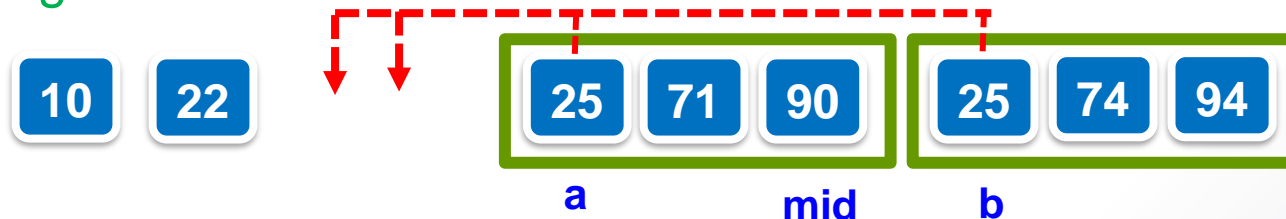
```
    }
```

```
} // end of while loop;
```

```
} // end of merge
```

# Merge (Pseudo Code)

```
else { // the 1st elements of the 2 halves are equal  
    if (they are the last elements) break;  
    1st element of 1st half joins end of the merged list;  
    move the 1st element of 2nd half to the end of the  
    merged list;  
}  
} // end of while loop;  
} // end of merge
```





# Merge (Pseudo Code)

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else { // the 1st elements of the 2 halves are equal
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}
} // end of while loop;
} // end of merge
```

**Challenge:**

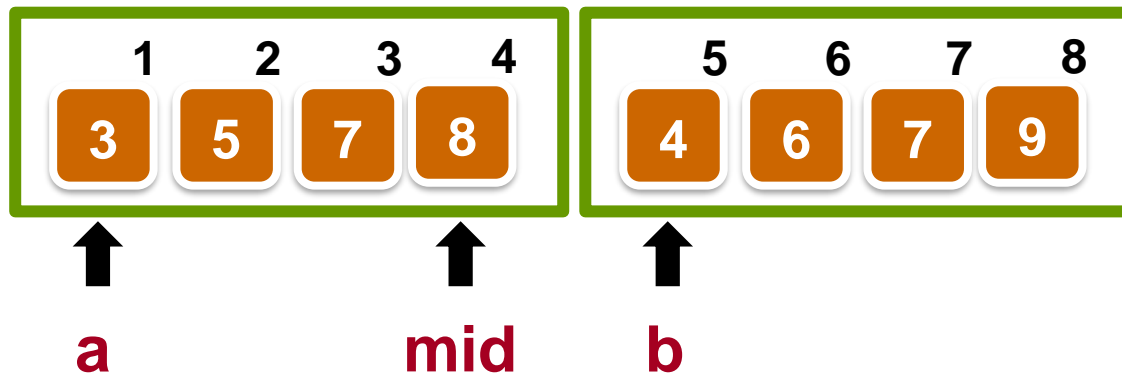
**How to do it without auxiliary  
storage for the merged list?**



# **Merge (Case Scenarios)**

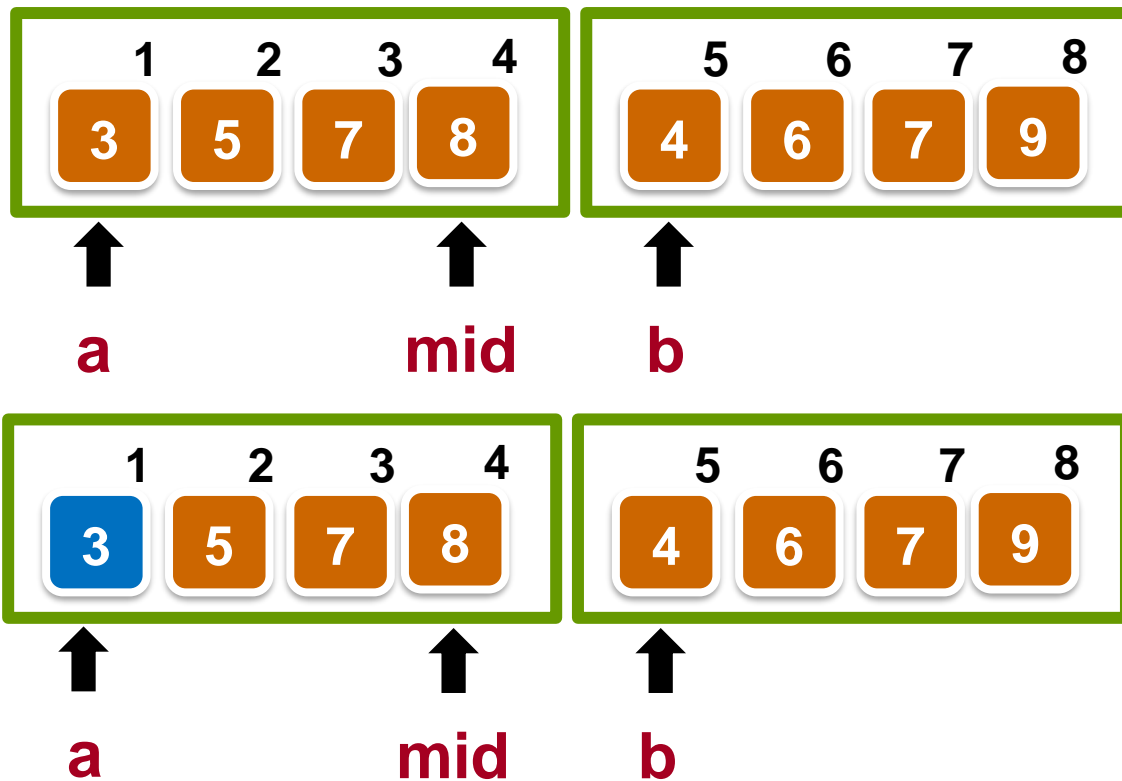
# Merge (Case Scenarios)

**Case 1:** 1st element of 1st half is smaller



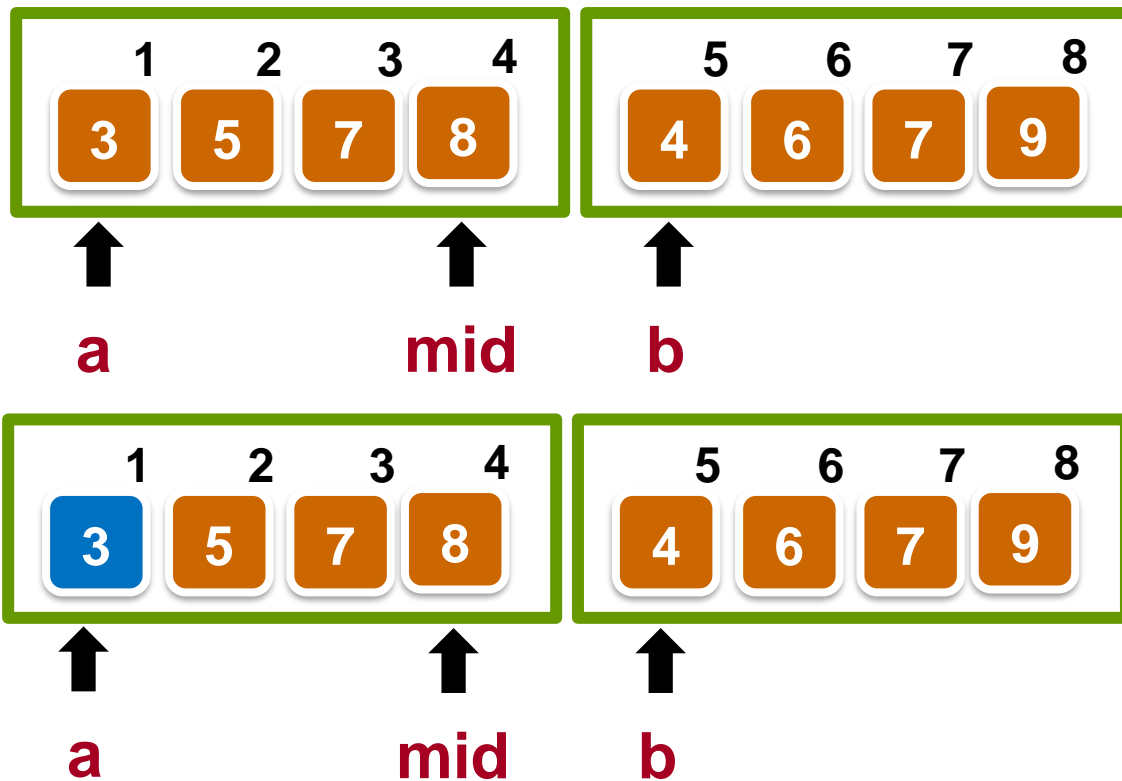
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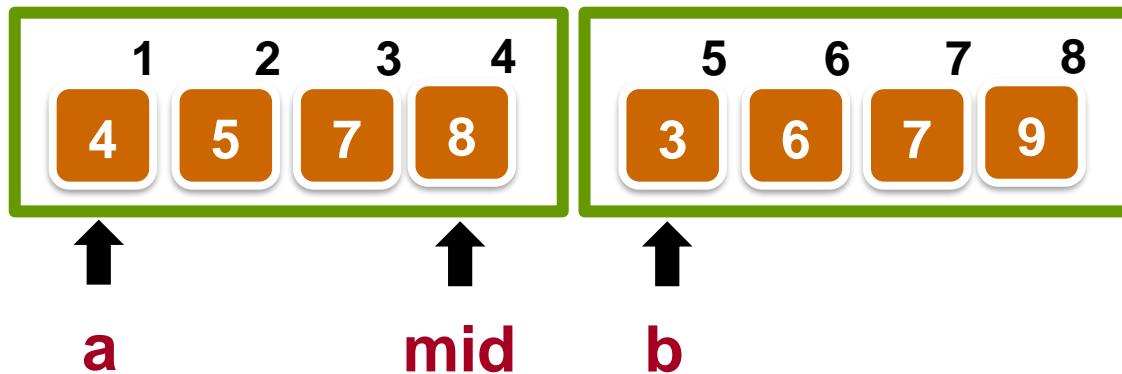
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**Case 1:** 1st element of 1st half is smaller



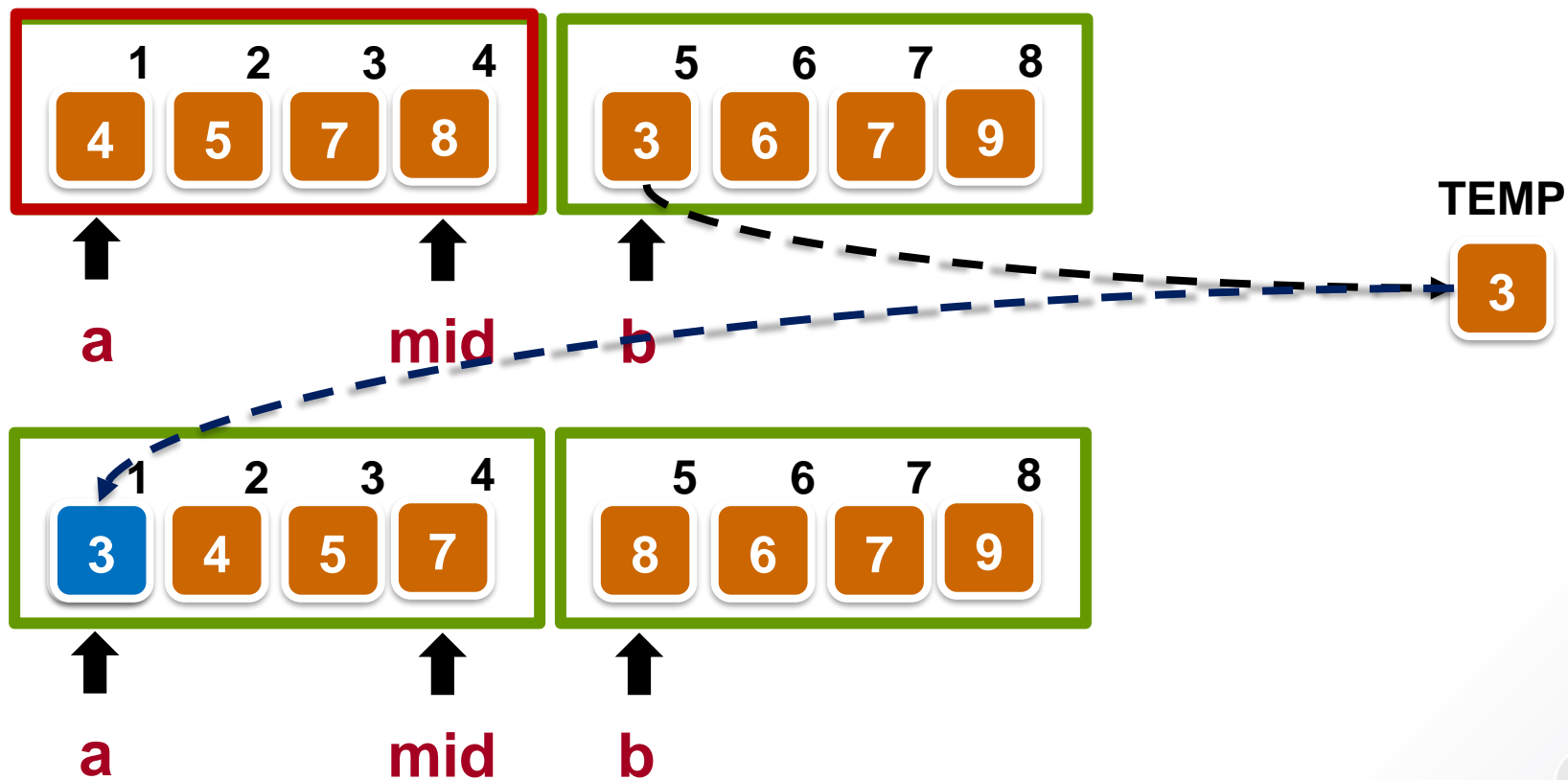
# Merge (Case Scenarios)

**Case 2:** 1st element of 2nd half is smaller



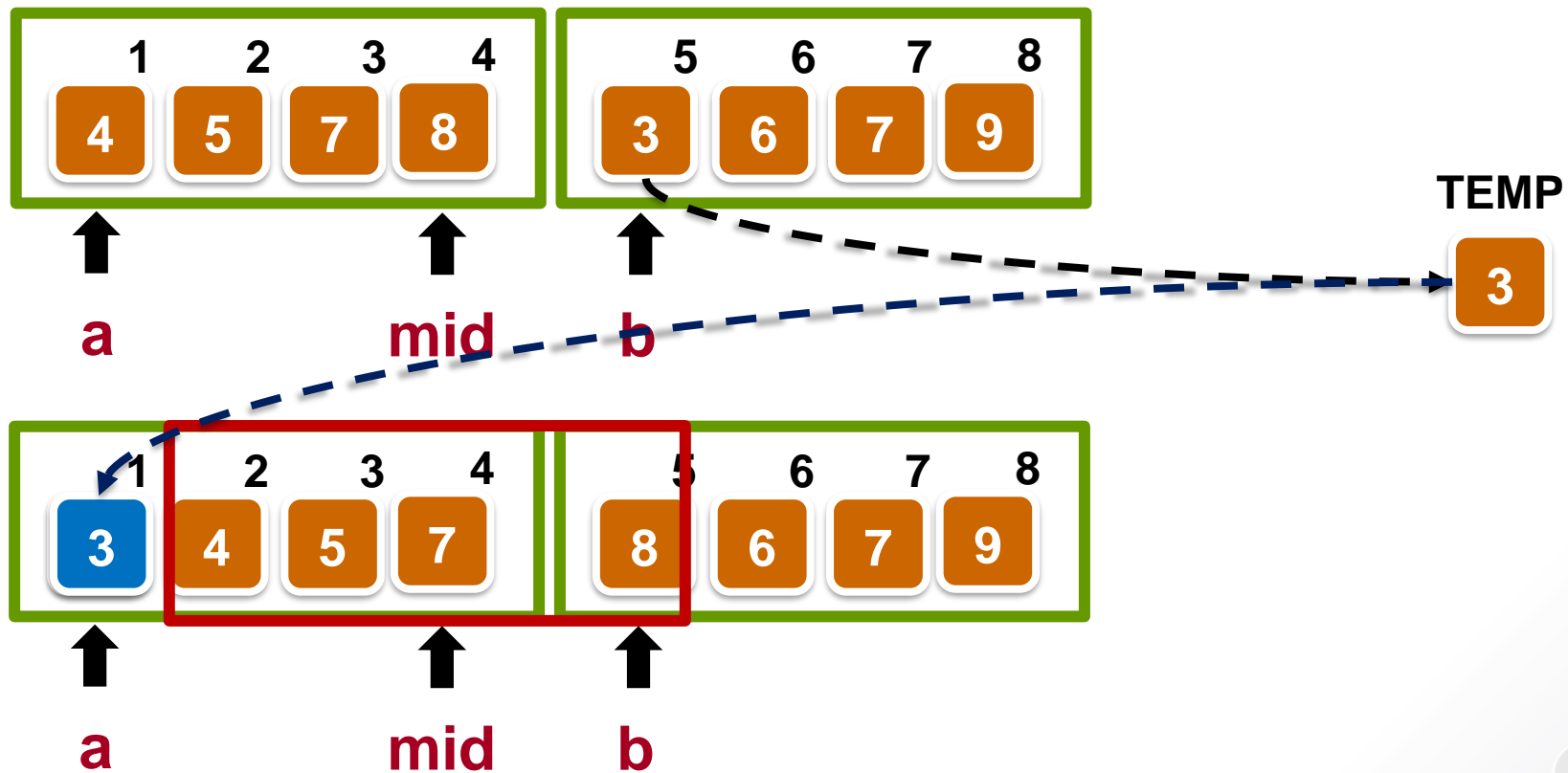
# Merge (Case Scenarios)

**Case 2:** 1st element of 2nd half is smaller



# Merge (Case Scenarios)

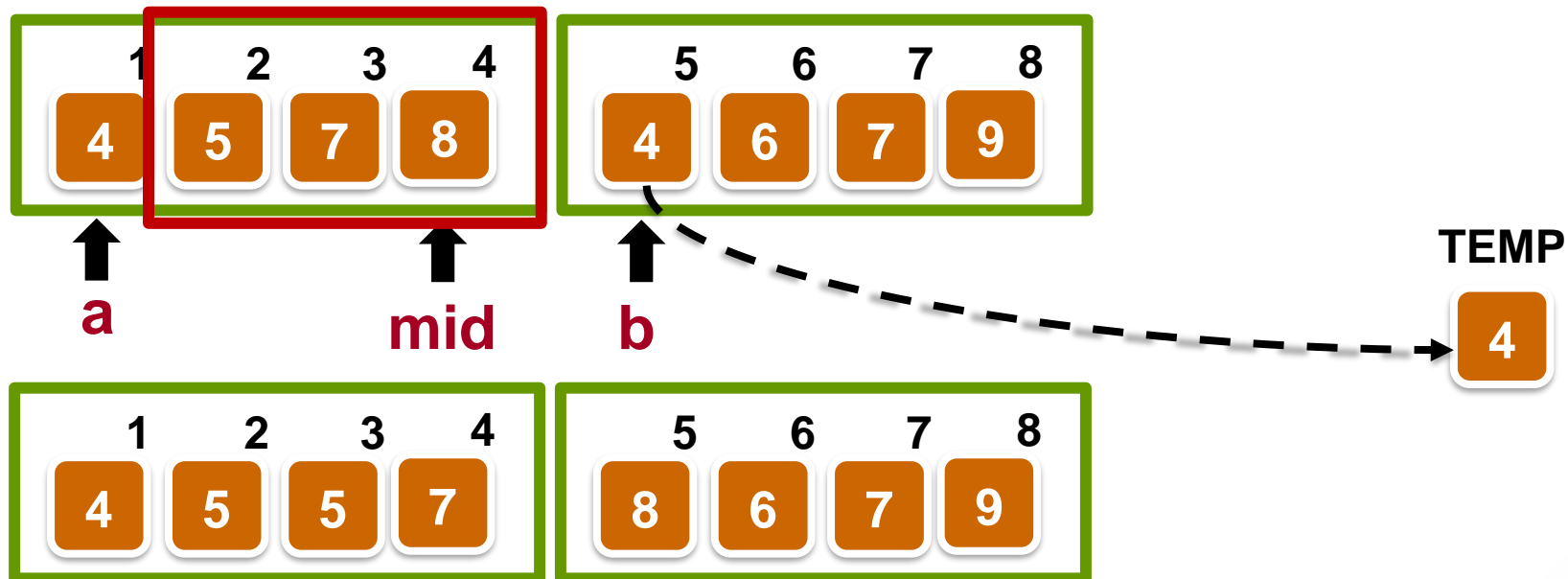
**Case 2:** 1st element of 2nd half is smaller





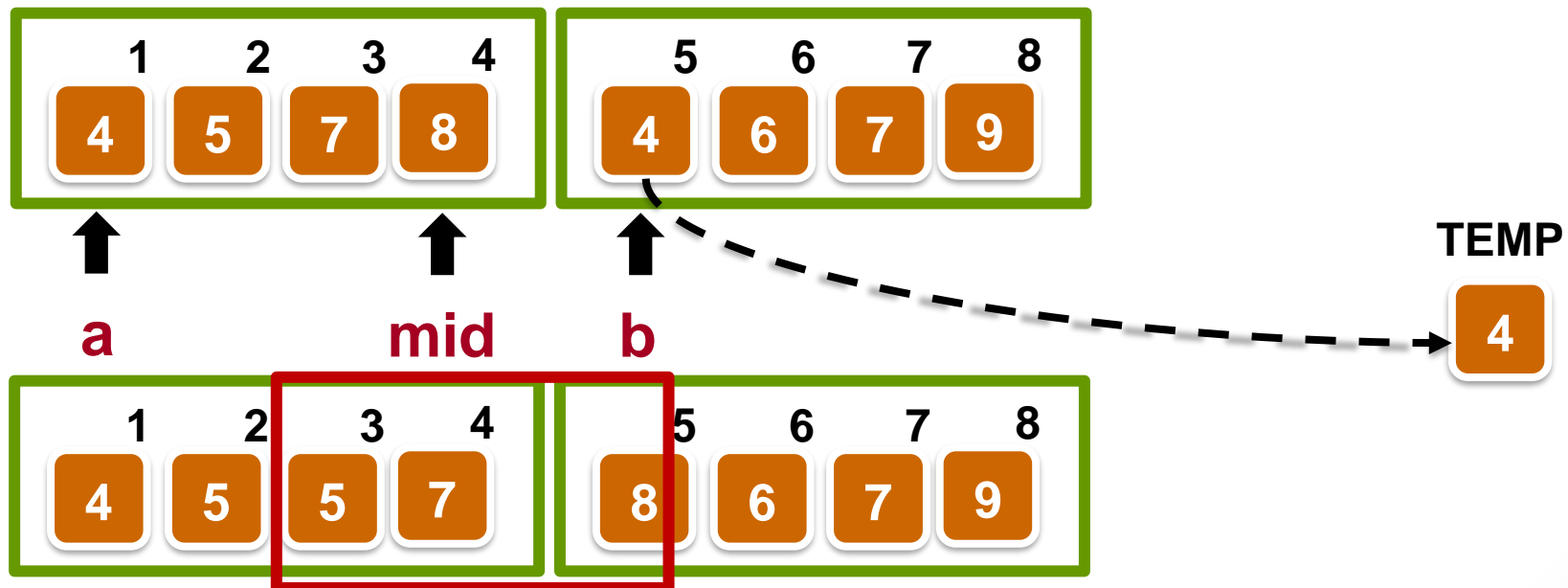
# Merge (Case Scenarios)

**Case 3:** 1st element of 2nd half is equal



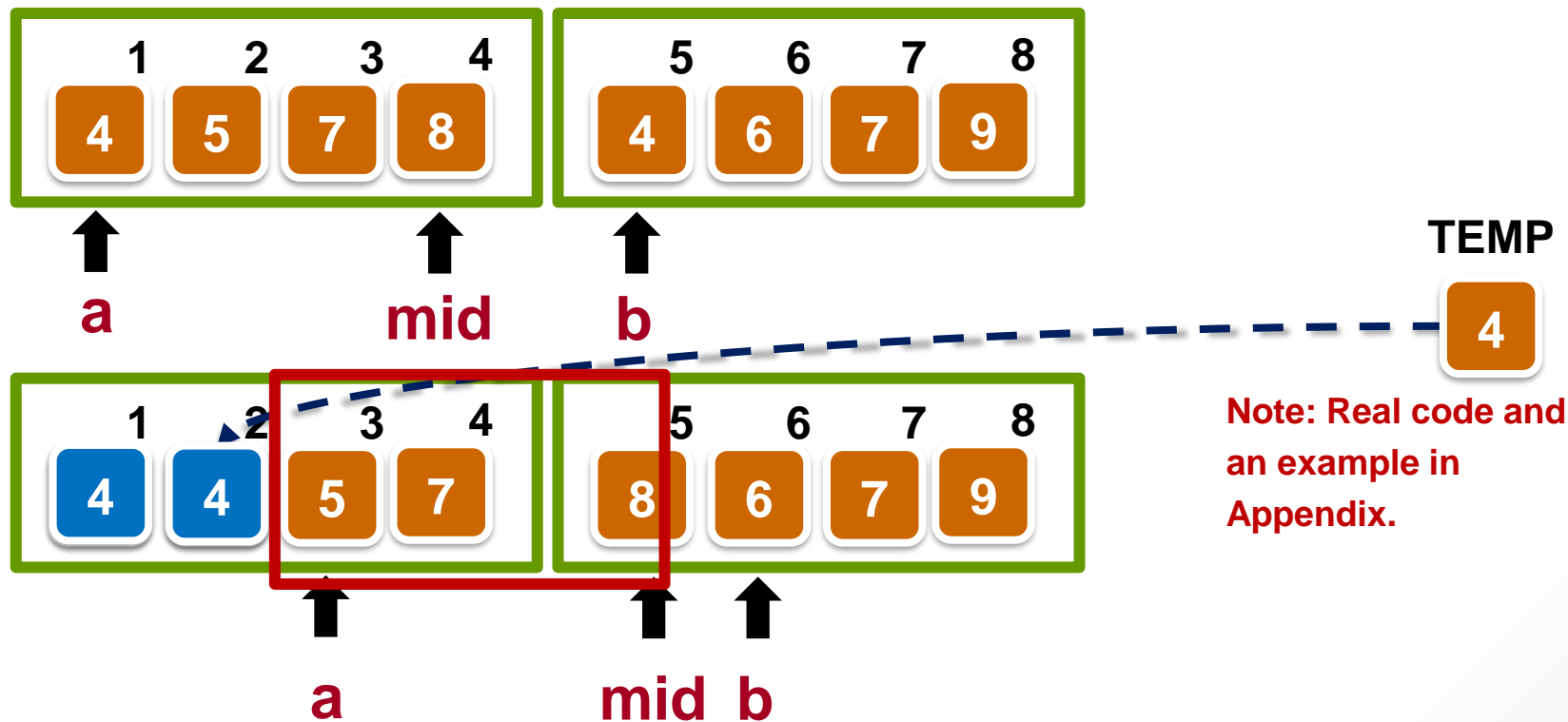
# Merge (Case Scenarios)

**Case 3:** 1st element of 2nd half is equal



# Merge (Case Scenarios)

**Case 3:** 1st element of 2nd half is equal



Note: Real code and an example in Appendix.

Can we just copy the first 4 twice?



# Mergesort Algorithm (Recap)

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- Since merging is performed directly on the original array, swapping and shifting are needed
- `mergesort()` partitions a contiguous array of elements between index `n` and `m` into two subarrays

```
void mergesort(int n, int m)
{
    int mid = (n+m)/2;
    if (m-n <= 0)
        return;
    else if (m-n > 1) {
        mergesort(n, mid);
        mergesort(mid+1, m);
    }
    merge(n, m);
}
```

# Mergesort Algorithm (Recap)

- Since merging is performed directly on the original array, swapping and shifting are needed
- `mergesort()` partitions a contiguous array of elements between index `n` and `m` into two subarrays
- Recursively partitions until `m-n ≤ 0`, then merge the resulting two subarrays

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void mergesort(int n, int m)
{   int mid = (n+m)/2;
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        return;
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# Mergesort Algorithm (Recap)

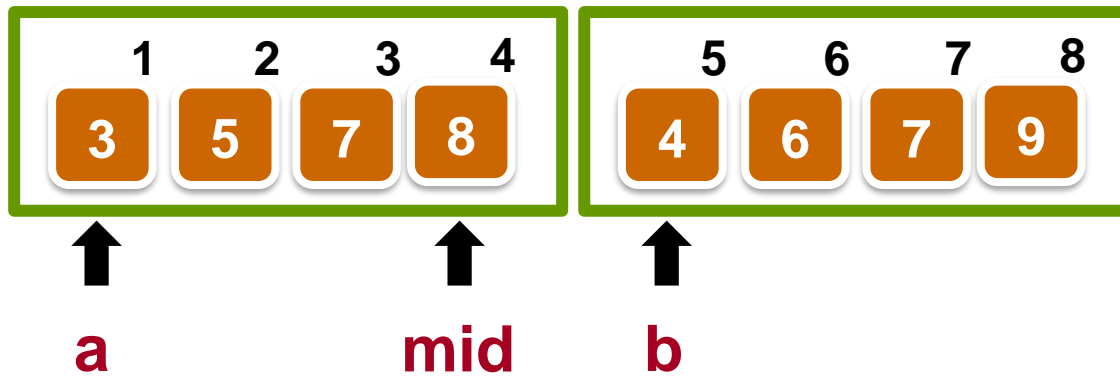
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- `merge()` function merges two sub-arrays of elements between index `n` and '`mid`', and between '`mid+1`' and `m`
- During merging, one element from each subarray is compared and the smaller one is inserted into new list

```
void mergesort(int n, int m)
{
    .....
    merge(n, m);
}
```



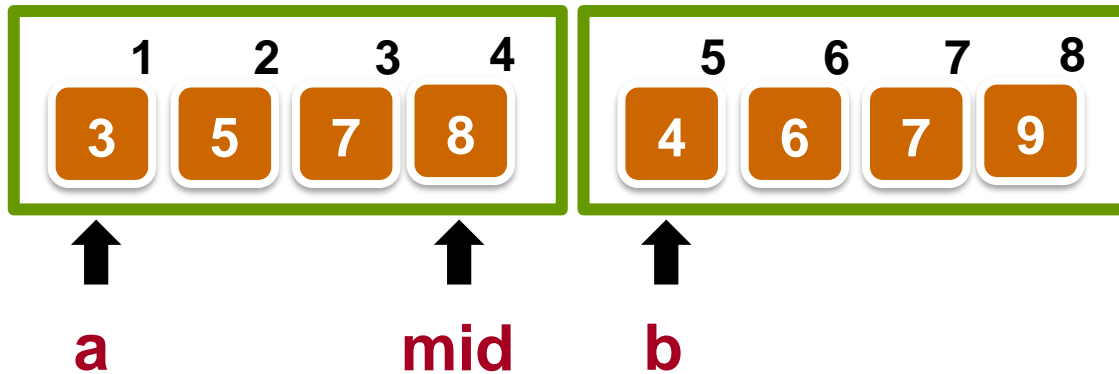
# Mergesort Algorithm (Recap)

- Left subarray runs from  $n$  to ' $mid$ ' with  $a$  as running index;  
right subarray runs from  $mid+1$  to  $m$  with  $b$  as running index



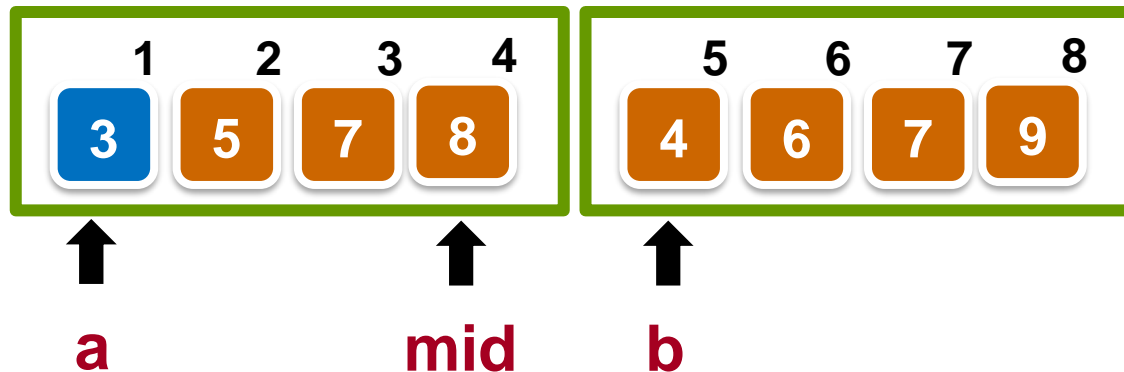
# Mergesort Algorithm (Recap)

- Left subarray runs from  $n$  to 'mid' with  $a$  as running index; right subarray runs from  $\text{mid}+1$  to  $m$  with  $b$  as running index
- $\text{slot}[a]$  is the head element of left subarray,  $\text{slot}[b]$  is the head element of right subarray



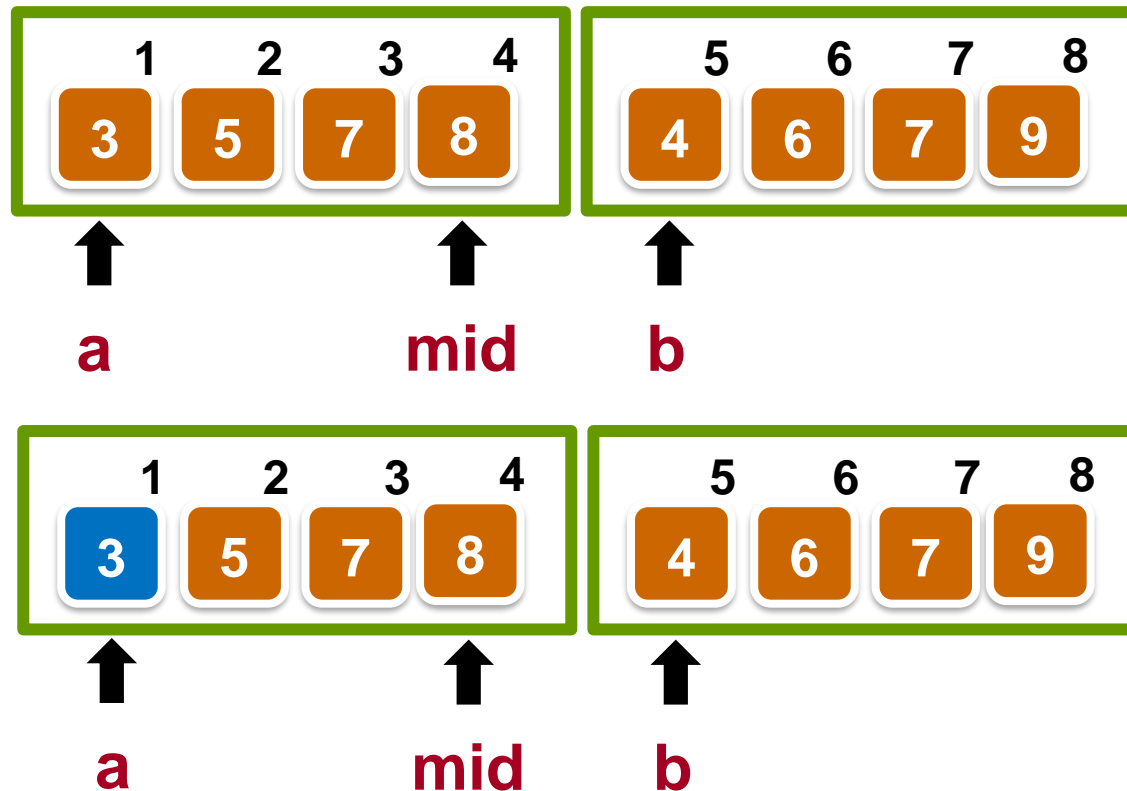
# Mergesort Algorithm (Recap)

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- $\text{slot}[a]$  is the head element of left subarray,  $\text{slot}[b]$  is the head element of right subarray
- During merging, both left and right subarrays shrink towards the right to make space for the newly merged array



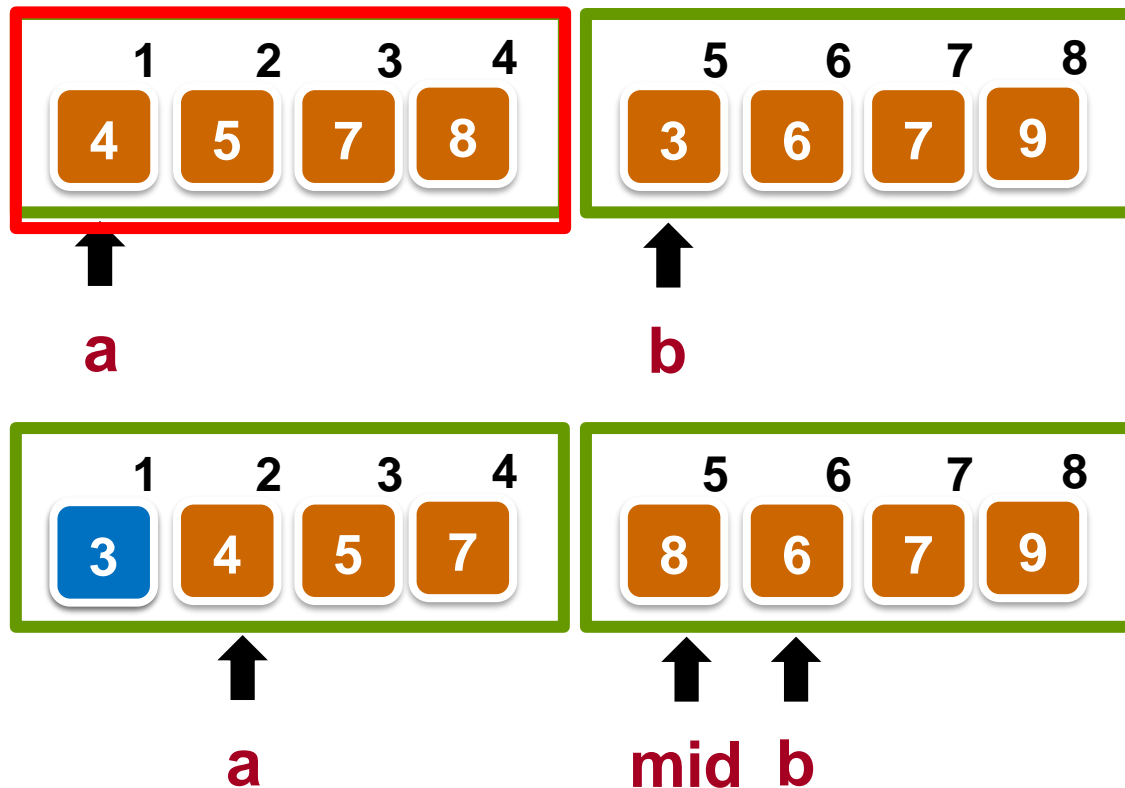
# Mergesort Algorithm (Recap)

**Case 1:** if  $\text{slot}[a] < \text{slot}[b]$ , there is nothing much to do since smaller element already in correct position (with regard to the merged array)



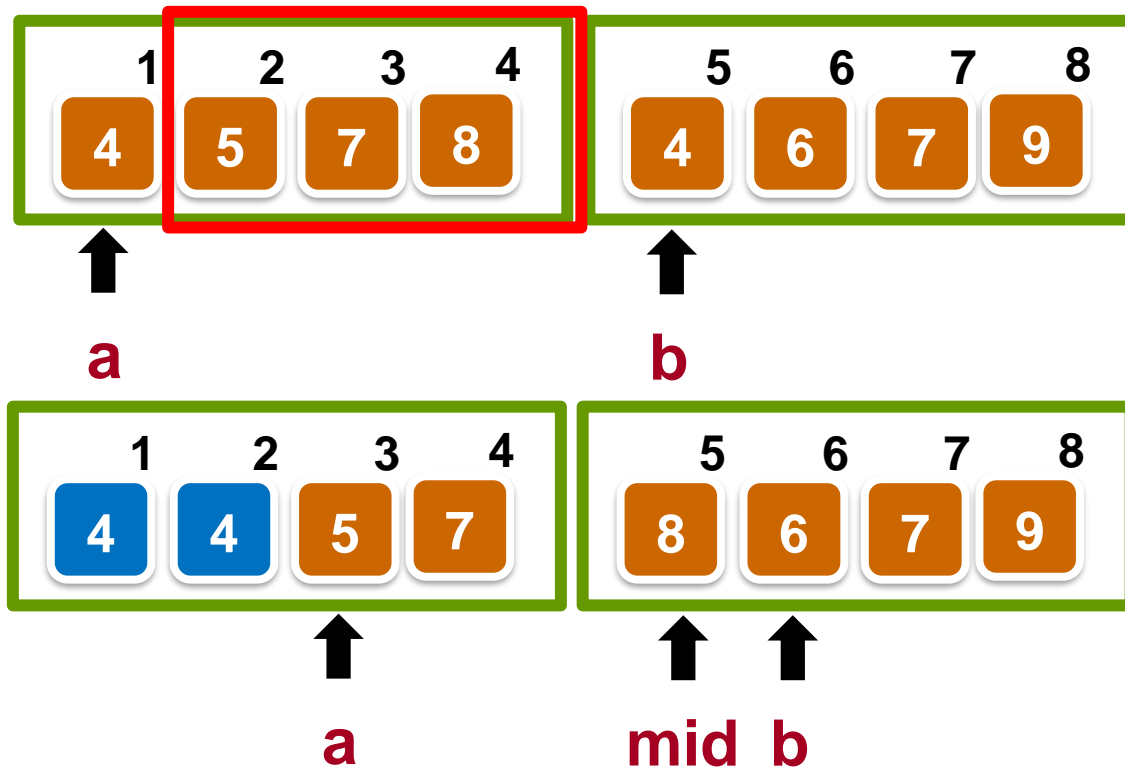
# Mergesort Algorithm (Recap)

**Case 2:** if  $\text{slot}[a] > \text{slot}[b]$ , then Right-shift (by one) elements of left subarray from index  $a$  to 'mid' and insert element at  $\text{slot}[b]$  into  $\text{slot}[a]$

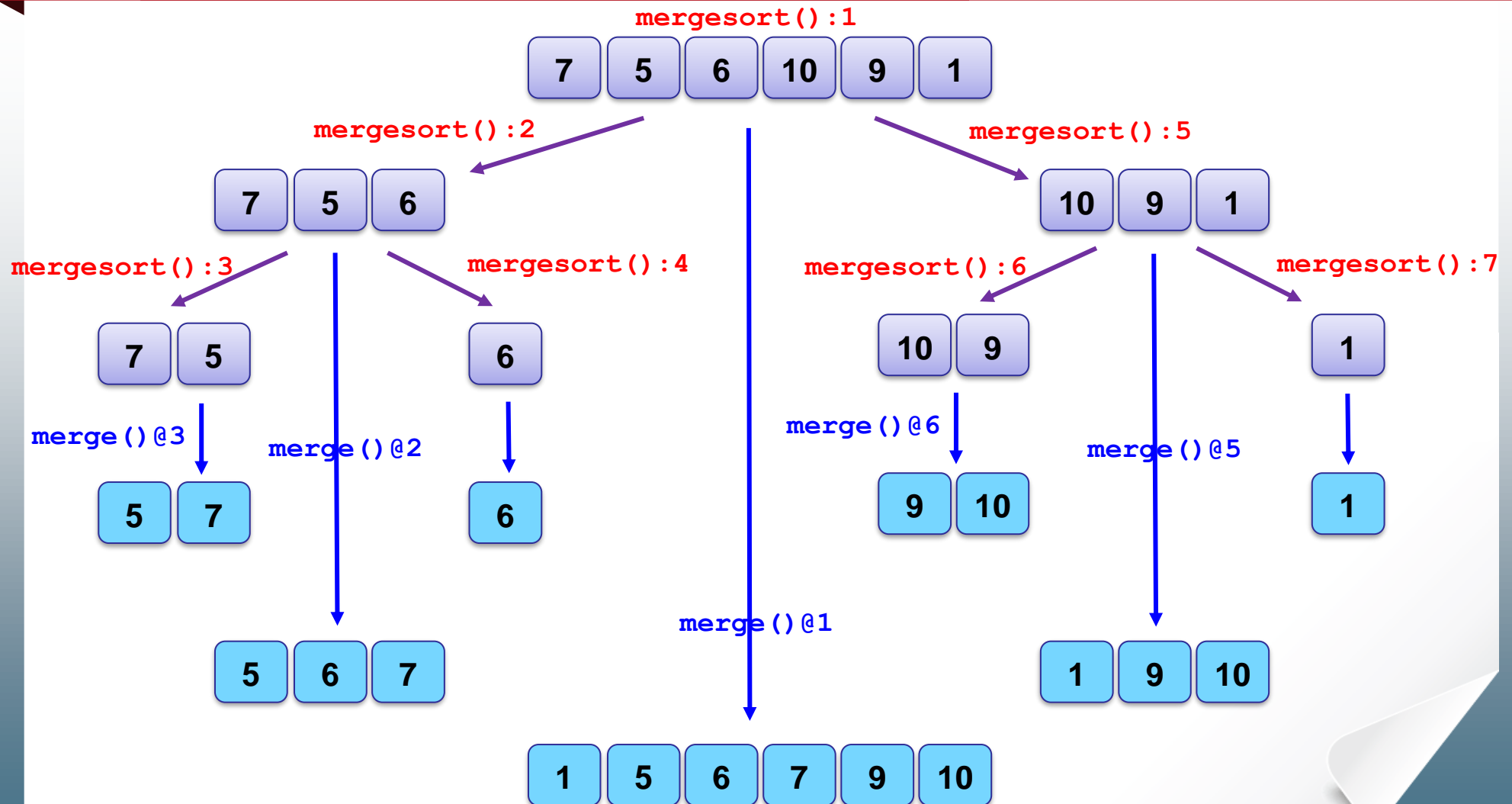


# Mergesort Algorithm (Recap)

**Case 3:** if  $\text{slot}[a] == \text{slot}[b]$ , then  $\text{slot}[a]$  is in the correct position. So, move  $\text{slot}[b]$  next to beside  $\text{slot}[a]$ , by Right-shifting and swapping



# Call Graph of Mergesort



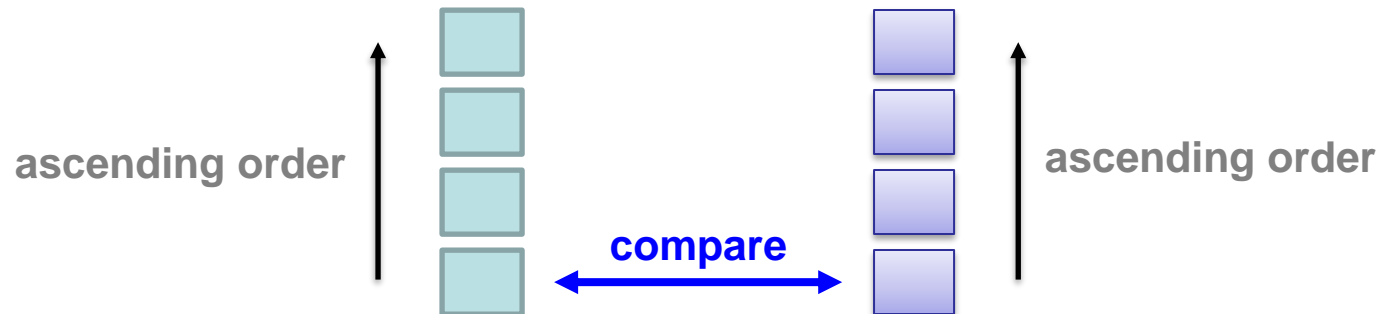


# Complexity of Mergesort



# Complexity of merge()

- After **each** comparison of keys from the two sub-lists, **at least one** element is moved to the new merged list and never compared again



- After the **last** key comparison, at least **two** elements will be moved into the merged list
- Thus, to merge two sub-lists of  $n$  elements in total, the number of key comparisons needed is **at most (worst case)**  $n - 1$ 
  - How about best case?

# Complexity of Mergesort

```
void mergesort(int s, int e) // s=start, e=end
```

```
{  int mid = (s+e)/2;
```

```
    if (e-s <= 0) return;
```

→  $W(1) = 0$

```
    else if (e-s > 1) {
```

```
        mergesort(s, mid);
```

→  $W(n/2)$

```
        mergesort(mid+1, e);
```

→  $W(n/2)$

$W(n)$

```
    }
```

```
    merge(s, e);
```

→ Worst case:  $n-1$

```
}
```

# Complexity of Mergesort

## Mergesort performance (assume $n = 2^k$ )

**Worst case :**

$$W(1) = 0,$$

$$W(n) = W(n/2) + W(n/2) + n-1 \quad \text{Or}$$

$$\begin{aligned} W(2^k) &= 2W(2^{k-1}) + 2^k - 1 \\ &= 2(2W(2^{k-2}) + 2^{k-1} - 1) + 2^k - 1 \\ &= 2^2W(2^{k-2}) + 2^k - 2 + 2^k - 1 \\ &= 2^2(2W(2^{k-3}) + 2^{k-2} - 1) + 2^k - 2 + 2^k - 1 \\ &= 2^3W(2^{k-3}) + 2^k - 2^2 + 2^k - 2 + 2^k - 1 \end{aligned}$$

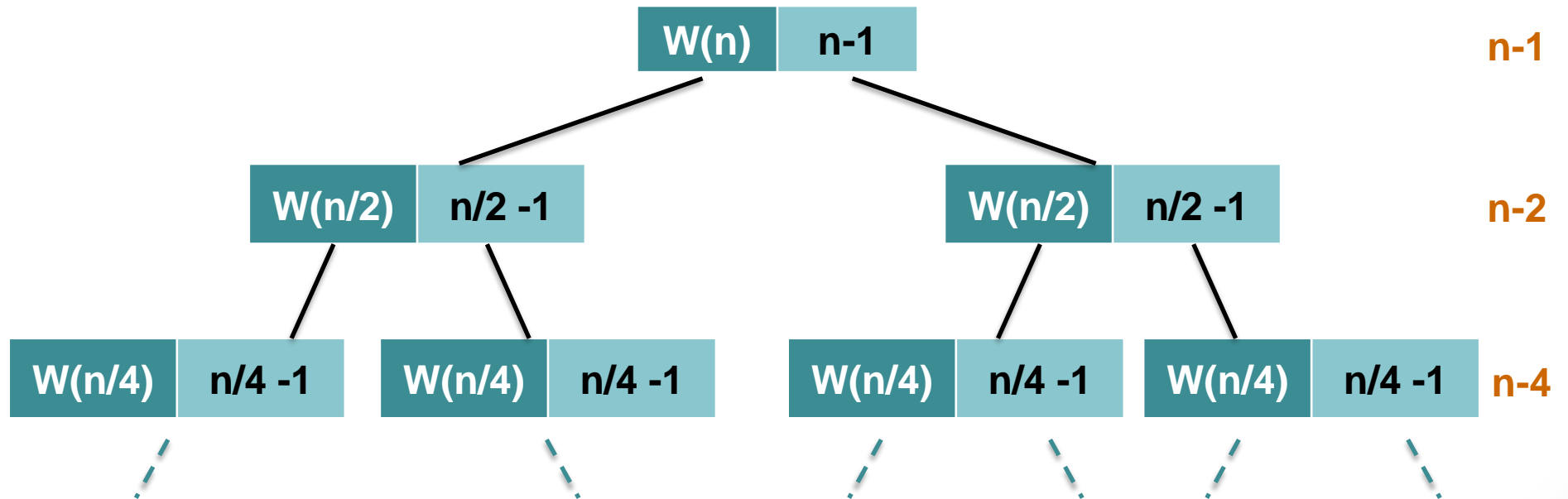
...

$$\begin{aligned} &= 2^kW(2^{k-k}) + k2^k - (1 + 2 + 4 + \dots + 2^{k-1}) \\ &= k2^k - (2^k - 1) \\ &= n \lg n - (n - 1) \\ &= O(n \lg n) \end{aligned}$$

$$k = \lg n$$

Geometric series

# Visually : Recursion Tree



$$W(2) = 2W(1) + 1 = 1$$

Total:  $O(n \lg n)$

Height of tree is  $k = O(\lg n)$

# Evaluation of Mergesort

## 😊 Strengths:

- 👉 Simple and good runtime behavior
- 👉 Easy to implement when using linked list

## 😞 Weaknesses:

- 👉 Difficult to implement for contiguous data storage such as array without auxiliary storage (requires data movements during merging)

# Summary

- Mergesort uses the **Divide and Conquer** approach.
  - It recursively divide a list into two halves of approximately equal sizes, until the sub-list is too small (no more than two elements).
  - Then, it recursively merges two sorted sub-lists into one sorted list.
- The worst-case running time for **merging** two sorted lists of total size  $n$  is  $n - 1$  key comparisons.
- The running time of Mergesort is  $O(n \lg n)$ .



# **SC2001/CE2101/CZ2101:**

# **Algorithm Design and Analysis**

## **Appendix**

## **(Merge operation in Mergesort)**

Instructor: Assoc. Prof. ZHANG Hanwang

# Merge Function

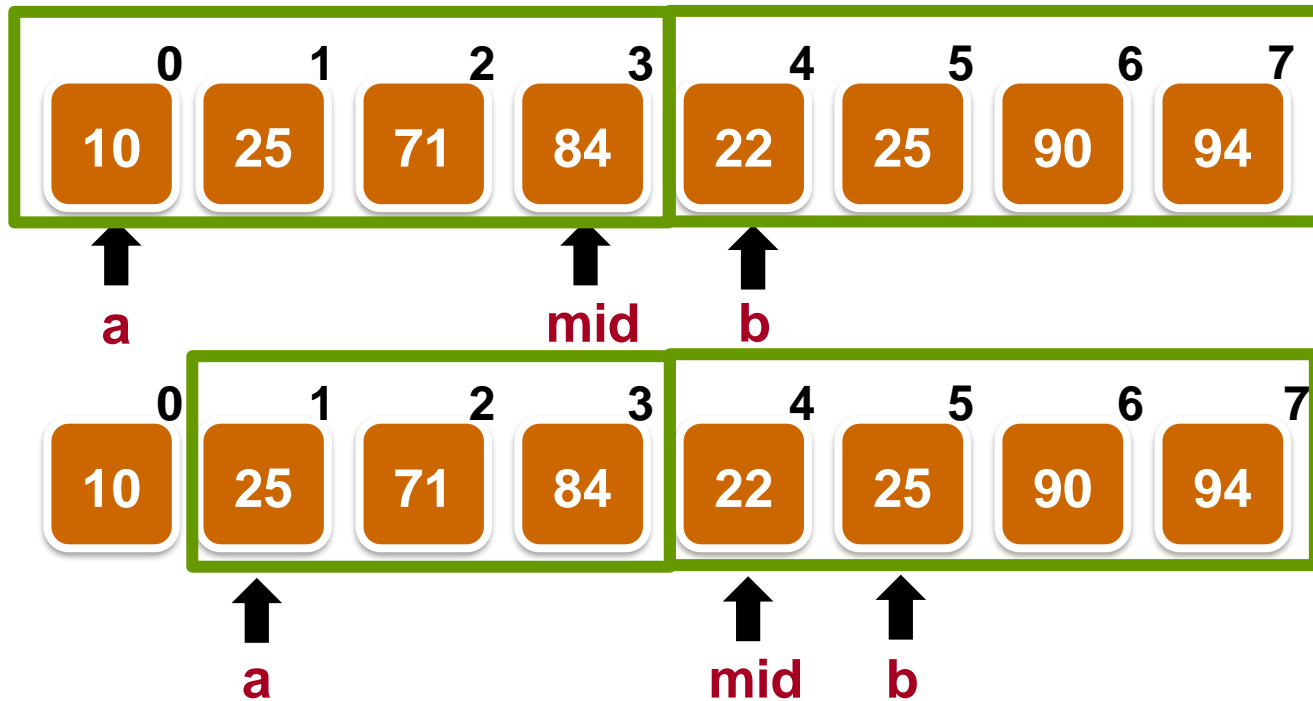
```
void merge(int n, int m)
{
    int mid = (n+m)/2;
    int a = n, b = mid+1, i, tmp;
    if (m-n <= 0) return;
    while (a <= mid && b <= m) {
        cmp = compare(slot[a], slot[b]);
        if (cmp > 0) { //slot[a] > slot[b]
            tmp = slot[b++];
            for (i = ++mid; i > a; i--)
                slot[i] = slot[i-1];
```



# Merge Function

```
        slot[a++] = tmp;
    } else if (cmp < 0) //slot[a] < slot[b]
        a++;
    else { //slot[a] == slot[b]
        if (a == mid && b == m)
            break;
        tmp = slot[b++];
        a++;
        for (i = ++mid; i > a; i--)
            slot[i] = slot[i-1];
        slot[a++] = tmp;
    }
} // end of while loop;
} // end of merge
```

# Merge Operation



## Parameters for merge:

$n:0, m:7$

$\text{mid} = (0+7)/2 = 3;$

$a = n; b = \text{mid}+1;$

## Comparison:

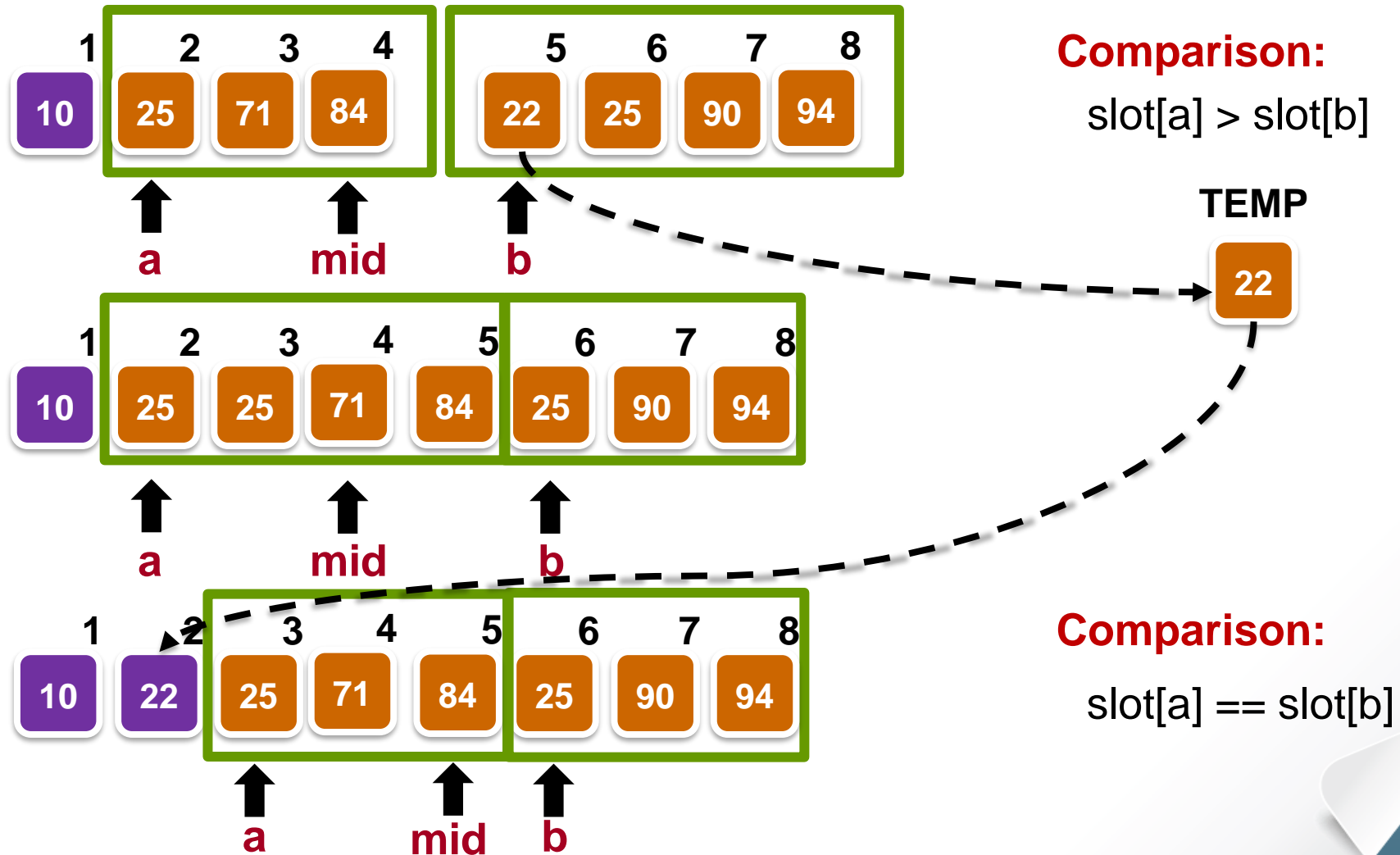
$\text{slot}[a] < \text{slot}[b]$

**a** : the 1<sup>st</sup> element of the 1<sup>st</sup> half

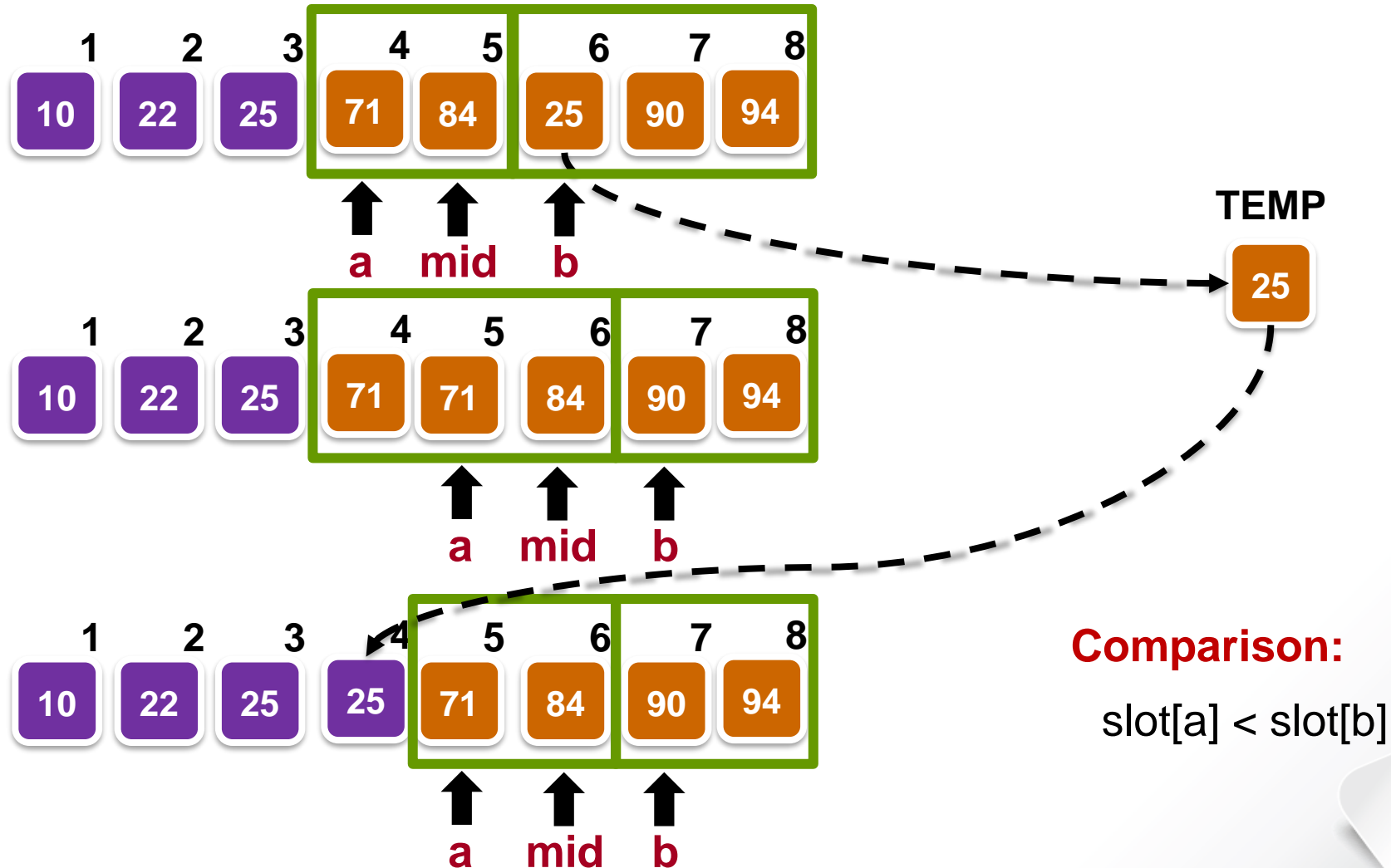
**mid** : the last element of the 1<sup>st</sup> half

**b** : the 1<sup>st</sup> element of the 2<sup>nd</sup> half

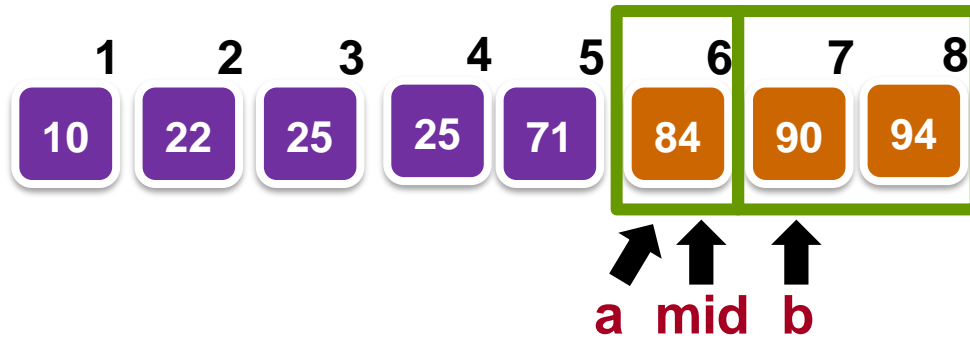
# Merge Operation



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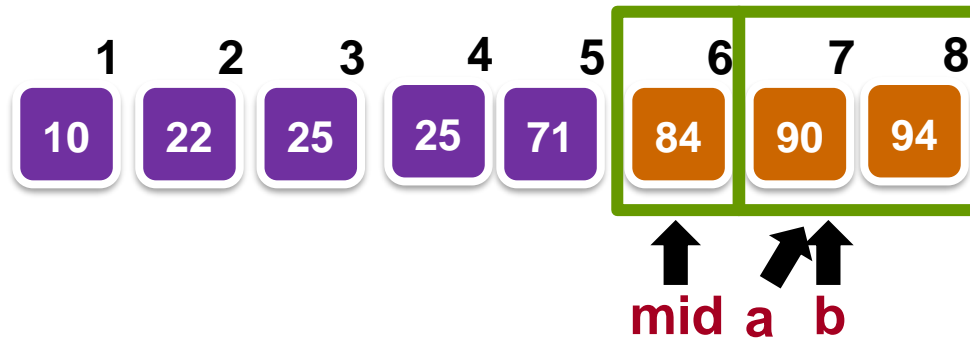


# Merge Operation



**Comparison:**

$\text{slot}[a] < \text{slot}[b]$



1<sup>st</sup> half empty

**Merge operation completed**