

696. Count Binary Substrings


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Approach #1: Group By Character [Accepted]

Intuition

We can convert the string `s` into an array `groups` that represents the length of same-character contiguous blocks within the string. For example, if `s = "110001111000000"`, then `groups = [2, 3, 4, 6]`.

For every binary string of the form `'0' * k + '1' * k` or `'1' * k + '0' * k`, the middle of this string must occur between two groups.

Let's try to count the number of valid binary strings between `groups[i]` and `groups[i+1]`. If we have `groups[i] = 2`, `groups[i+1] = 3`, then it represents either `"0011"` or `"1100"`. We clearly can make `min(groups[i], groups[i+1])` valid binary strings within this string. Because the binary digits to the left or right of this string must change at the boundary, our answer can never be larger.

Algorithm

Let's create `groups` as defined above. The first element of `s` belongs in it's own group. From then on, each element either doesn't match the previous element, so that it starts a new group of size 1; or it does match, so that the size of the most recent group increases by 1.

Afterwards, we will take the sum of `min(groups[i-1], groups[i])`.

Python

```
class Solution(object):
    def countBinarySubstrings(self, s):
        groups = [1]
        for i in xrange(1, len(s)):
            if s[i-1] != s[i]:
                groups.append(1)
            else:
                groups[-1] += 1

        ans = 0
        for i in xrange(1, len(groups)):
            ans += min(groups[i-1], groups[i])
        return ans
```

Alternate Implentation

```
class Solution(object):
    def countBinarySubstrings(self, s):
        groups = [len(list(v)) for _, v in itertools.groupby(s)]
        return sum(min(a, b) for a, b in zip(groups, groups[1:]))
```

Java

```

class Solution {
    public int countBinarySubstrings(String s) {
        int[] groups = new int[s.length()];
        int t = 0;
        groups[0] = 1;
        for (int i = 1; i < s.length(); i++) {
            if (s.charAt(i-1) != s.charAt(i)) {
                groups[++t] = 1;
            } else {
                groups[t]++;
            }
        }

        int ans = 0;
        for (int i = 1; i <= t; i++) {
            ans += Math.min(groups[i-1], groups[i]);
        }
        return ans;
    }
}

```

Complexity Analysis

- Time Complexity: $O(N)$, where N is the length of s . Every loop is through $O(N)$ items with $O(1)$ work inside the for-block.
- Space Complexity: $O(N)$, the space used by `groups`.

Approach #2: Linear Scan [Accepted]

Intuition and Algorithm

We can amend our *Approach #1* to calculate the answer on the fly. Instead of storing `groups`, we will remember only `prev = groups[-2]` and `cur = groups[-1]`. Then, the answer is the sum of `min(prev, cur)` over each different final `(prev, cur)` we see.

Python

```

class Solution(object):
    def countBinarySubstrings(self, s):
        ans, prev, cur = 0, 0, 1
        for i in xrange(1, len(s)):
            if s[i-1] != s[i]:
                ans += min(prev, cur)
                prev, cur = cur, 1
            else:
                cur += 1

        return ans + min(prev, cur)

```

Java

```

class Solution {
    public int countBinarySubstrings(String s) {
        int ans = 0, prev = 0, cur = 1;
        for (int i = 1; i < s.length(); i++) {
            if (s.charAt(i-1) != s.charAt(i)) {
                ans += Math.min(prev, cur);
                prev = cur;
                cur = 1;
            } else {
                cur++;
            }
        }
        return ans + Math.min(prev, cur);
    }
}

```

Complexity Analysis

- Time Complexity: $O(N)$, where N is the length of s . Every loop is through $O(N)$ items with $O(1)$ work inside the for-block.
- Space Complexity: $O(1)$, the space used by `prev`, `cur`, and `ans`.

Analysis written by: @awice (<https://leetcode.com/awice>).

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makikvues (/makikvues) ★ 3 ⌚ June 30, 2018 8:41 AM

```

public int CountBinarySubstrings(string s)
{
    int count = 0;
    int consecutiveCount = 1;

```

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dbalagula (/dbalagula) ★ 1 ⌚ March 23, 2018 11:53 PM

```

var i = 0;
var sum = 0;
while (i < s.length){
    var type = (s[i] == 0 ? 0 : 1);
    if (type == 0){

```

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can't understand, it's too hard for a rookie

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