

Probability and Recurrences

Reminders and Leftovers

- Assignment 7 is due tonight
 - Help hours today: 1.30-3pm (me), 3-5 pm and 5-11pm (TAs)
- Assignment 8 will be released today; due next Wed
- Where we are:
 - Last lecture we introduced basics of probability (sample space, events, independence, conditional probability)
 - Saw some examples
 - Today: we'll define random variable, expectation and see examples of analyzing expectation of probabilistic processes

The Birthday Paradox

- Suppose that there are m students in a lecture hall
- Assume for each student, any of the $n = 365$ possible days are equally likely as their birthday
- Assume birthday are mutually independent
- **Question.** What is the likelihood that no two students have the same birthday?
- Let A_i be the event that the i th persons birthday is different from the previous $i - 1$ people
- \Pr (all m different birthdays)
= $\Pr(A_1 \cap A_2 \cap \dots \cap A_m)$
= $\Pr(A_1) \cdot \Pr(A_2 | A_1) \cdot \Pr(A_3 | A_1 \cap A_2) \dots \Pr(A_n | A_1 \cap \dots \cap A_{n-1})$



The Birthday Paradox

- Pr (all m different birthdays)

$$= 1 \cdot \left(1 - \frac{1}{n}\right) \cdot \left(1 - \frac{2}{n}\right) \cdot \left(1 - \frac{3}{n}\right) \dots \left(1 - \frac{m-1}{n}\right)$$

$$= \prod_{j=1}^{m-1} \left(1 - \frac{j}{n}\right) \leq \prod_{j=1}^{m-1} e^{-j/n} = e^{-1/n(\sum_{j=1}^{m-1} j)} \approx e^{-m^2/2n}$$

- $m \approx \sqrt{2n \ln 2}$ for probability to be 1/2
- For $n = 365$, we get $m = 22.49$
- Thus, with around 23 people in this class, we have a 50% chance of two people having the same birthday

Important Inequality:

$$(1 - x) \leq \left(\frac{1}{e}\right)^x \text{ for } x \geq 1$$

Random Variable

- Event either does or does not happen, what if we want to capture *magnitude* of a probabilistic event
- Suppose I flip n independent fair coins: the # of heads is a random variable
- Number that comes up when we roll a fair die is a random variable
- If an algorithm flips some coins then the running time of the algorithm is a random variable
- **Definition.** A random variable X is a function from a sample space \mathcal{S} (with a probability measure) to some value set (e.g. real numbers, integers, etc.)

Random Variable: Example

- So for example I flip a coin 10 times. Let X be the number of heads
 - $\Pr[X = 0] = 1/2^{10}$
 - $\Pr[X = 10] = 1/2^{10}$
 - $\Pr[X = 4] ?$
 - $\Pr[X = 4] = \binom{10}{4} \frac{1}{2^4} \frac{1}{2^6} = \frac{105}{512}$
- A random variable that is 0 or 1 (indicating if something happens or not) is called an *indicator random variable or Bernoulli random variable*

Expectation

- Every time you do the experiment, associated random variable takes a different value
- How can we characterize the average behavior of a random variable?
- **Definition.** Expected value of a random variable R defined on a sample space S is

$$E(R) = \sum_{w \in S} R(w) \cdot \Pr(w)$$

- Let R be the number that comes up when we roll a fair, six-sided die, then the expected value of R is

$$E(R) = \sum_{i=1}^6 i \cdot \frac{1}{6} = \frac{1}{6}(1 + 2 + 3 + 4 + 5 + 6) = \frac{7}{2}$$

To get the E to look good in latex,
use `\mathrm{E}`

(We won't use it like \mathbb{E} in this class, but
if you really want to, it's `\mathbb{E}`)

Expectation

- We can group together outcomes for which the random variable takes the same value
- **Alternate Definition.** Expected value of a random variable R defined on a sample space S is

$$E(R) = \sum_x x \cdot \Pr(R = x)$$

- If A is an arbitrary event with $\Pr[A] > 0$, the conditional expectation of X given A is

$$E[X | A] := \sum_x x \cdot \Pr[X = x | A]$$

- **(Law of total expectation)** If $\{A_1, A_2, \dots\}$ is a finite partition of the sample space:

$$E(X) = \sum_i E(X | A_i) \cdot \Pr(A_i)$$



Very useful !

Linearity of Expectation

- *Very important* tool in randomized algorithm
- Expectation of random variables obey a wonderful rule
- Informally, it says that the expectation of a sum is the sum of the expectations.
- Formally, for any random variables X_1, X_2, \dots, X_n and any coefficients $\alpha_1, \alpha_2, \dots, \alpha_n$

$$\mathbb{E}\left[\sum_{i=1}^n (\alpha_i \cdot X_i)\right] = \sum_{i=1}^n (\alpha_i \cdot \mathbb{E}[X_i])$$

Very useful !

- **Note.** Always true! Linearity of expectation **does not require independence** of random variables.

Bernoulli Distribution

- A probability distribution assigns a probability to each possible value of a random variable
- Suppose you run an experiment with probability of success p and failure $1 - p$.
- Example, coin toss where head is success.
- Let X be a Bernoulli or indicator random variable that is **1** if we succeed, and **0** otherwise. Then,

$$E[X] = \sum_x x \cdot \Pr[X = x] = 0 \cdot \Pr[X = 0] + 1 \cdot \Pr[X = 1] = p$$

- **Remember this:** expectation of an indicator random variable is exactly the probability of success!



Expected Success: n Bernoulli Trials

- Consider n independent Bernoulli trials (with success probability p). Let R denote the number of successes
 - R is said to follow a *Binomial distribution* (we'll revisit this)
- We want to know expected number of successes $E(R)$
- Can write R as a sum of indicator random variables

$$R = \sum_i R_i \text{ where } R_i = 0 \text{ or } R_i = 1$$

- Then $E[R] = E \left[\sum_i R_i \right]$, how can we simplify this by LoE?

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- Can write R as a sum of indicator random variables

- $R = \sum_i R_i$ where $R_i = 0$ or $R_i = 1$

- Then $E[R] = E \left[\sum_i R_i \right] = \sum_i E[R_i] = \sum_{i=1}^n p = np$

Hat Check Problem

- There is a dinner party where n people check their hats. The hats get mixed up during dinner, so that afterward each person receives a random hat.
- What is the expected number of people who get their own hat?

Hat Check Problem

- There is a dinner party where n people check their hats. The hats get mixed up during dinner, so that afterward each person receives a random hat.
- What is the expected number of people who get their own hat?
- Let R be the random variable denoting the number of men who get their hat back. Goal: compute $E(R)$.
- Usual trick. Express random variable R as a sum of indicator random variables R_i is 1 if i th person gets their hat back, else it is 0.

- Then,
$$R = \sum_{i=1}^n R_i$$



Hat Check Problem

- What is $E(R_i)$?
- $R_i = 1$ if i gets his hat; $R_i = 0$ otherwise
- By definition,
$$E(R_i) = 1 \cdot \Pr(i \text{ gets } i\text{'s hat}) + 0 \cdot \Pr(i \text{ gets another hat})$$
- $E(R_i) = \Pr(i \text{ gets } i\text{'s hat})$
- Need $\Pr(i \text{ gets } i\text{'s hat})$
- Sample space: All orderings of hats $n!$ (returned in order)
- Number of outcomes where i 's gets correct hat?
 - $(n - 1)!$



Hat Check Problem

- What is $E(R_i)$?
- $R_i = 1$ if i gets his hat; $R_i = 0$ otherwise
- By definition,
$$E(R_i) = 1 \cdot \Pr(i \text{ gets } i\text{'s hat}) + 0 \cdot \Pr(i \text{ gets another hat})$$
- $E(R_i) = \Pr(i \text{ gets } i\text{'s hat}) = \frac{(n-1)!}{n!} = \frac{1}{n}$
- $$E(R) = E\left(\sum_{i=1}^n R_i\right) = \sum_{i=1}^n E(R_i) = \sum_{i=1}^n \frac{1}{n} = 1$$
- In expectation, one person gets their hat back!



Uniform Distribution

- When every outcome is equally likely
- Let X be the random variable of the experiment and S be the sample space

- $\Pr[X = x] = \frac{1}{|S|}$

- $E[X] = \frac{1}{|S|} \cdot \sum_{x \in S} \Pr(X = x)$

- Example
 - fair coin toss: heads and tails are equally likely
 - fair die roll: all numbers are equally likely



Card Guessing: Memoryless

- To entertain your family you have them shuffle deck of n cards and then turn over one card at a time. Before each card is turned, you predict its identity. You have no psychic abilities or memory to remember cards
- Your strategy: guess uniformly at random
- How many predictions do you expect to be correct?
- Let X denote the r.v. equal to the number of correct predictions and X_i denote the indicator variable that the i th guess is correct

- Thus, $X = \sum_{i=1}^n X_i$ and $E[X] = E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i]$

- $E[X_i] = 0 \cdot \Pr(X_i = 0) + 1 \cdot \Pr(X_i = 1) = \Pr(X_i = 1) = 1/n$

- Thus, $E[X] = 1$



Card Guessing: Memoryfull

- Suppose we play the same game but now assume you have the ability to remember cards that have already been turned
- Your strategy: guess uniformly at random among cards that have not been turned over
- Let X denote the r.v. equal to the number of correct predictions and X_i denote the indicator variable that the i th guess is correct

- Thus, $X = \sum_{i=1}^n X_i$ and $E[X] = E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i]$

- $E[X_i] = \Pr(X_i = 1) = \frac{1}{n - i + 1}$

- Thus, $E[X] = \sum_{i=1}^n \frac{1}{n - i + 1} = \sum_{i=1}^n \frac{1}{i}$

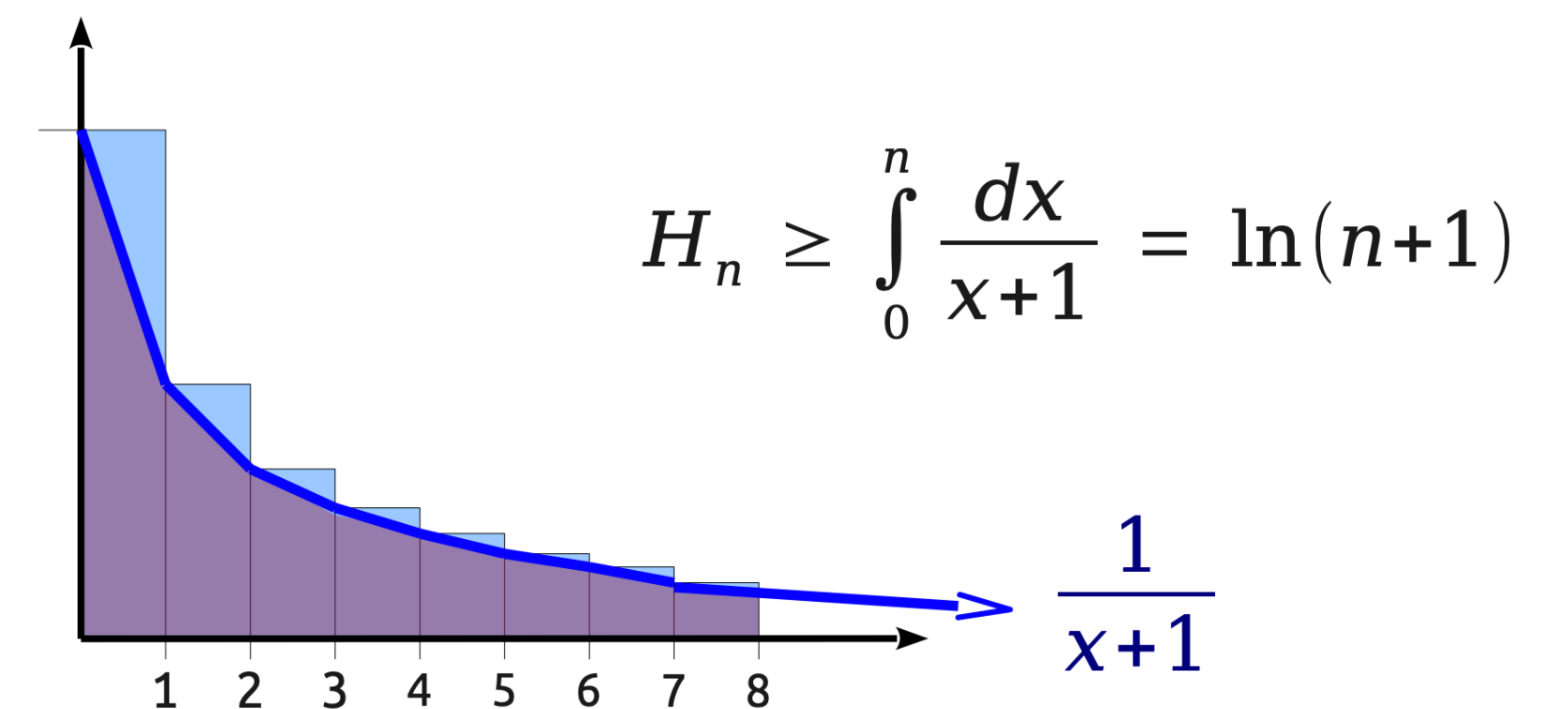
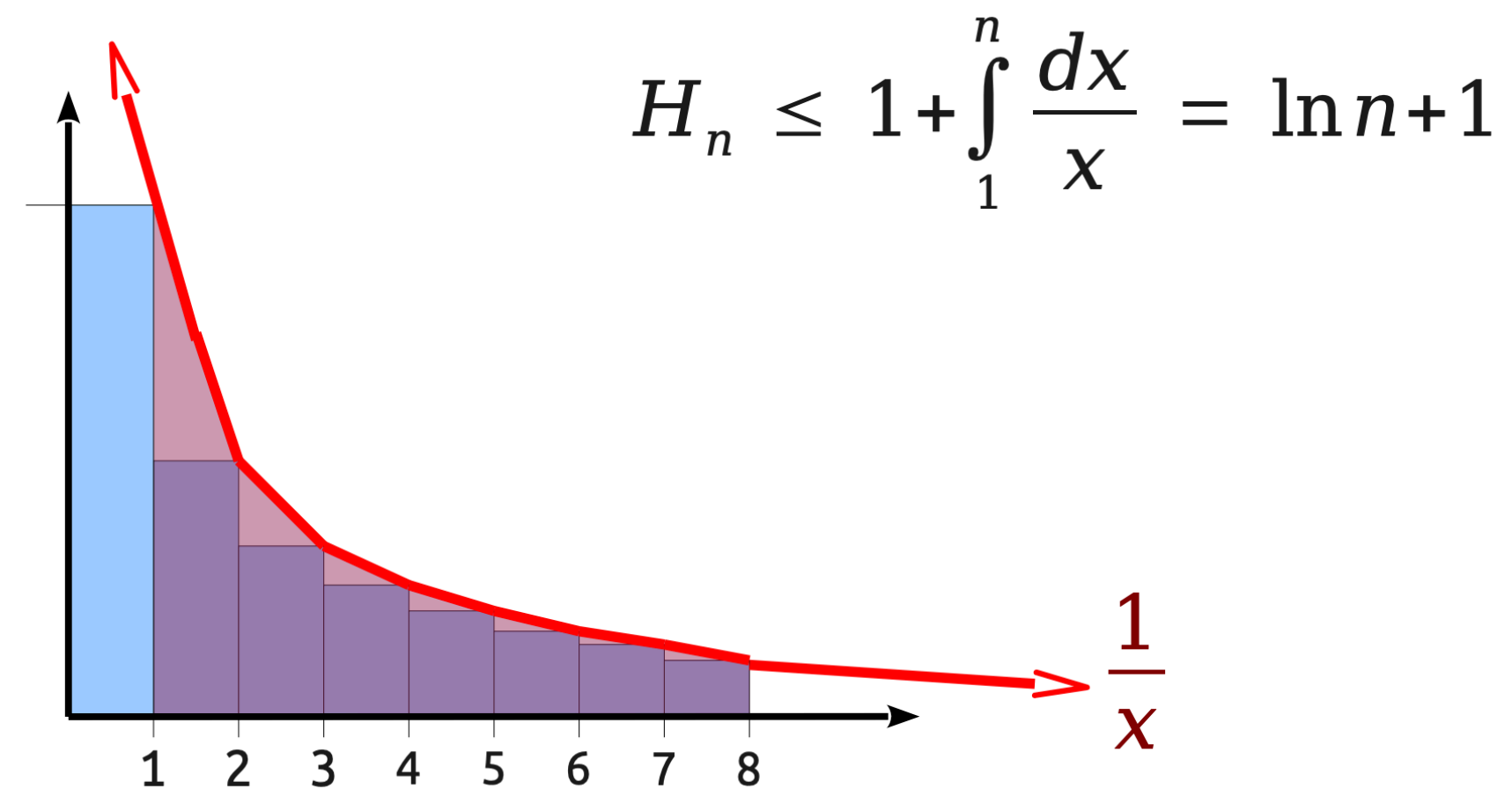


Harmonic Numbers

- The n th harmonic number, denoted H_n is defined as

$$H_n = \sum_{i=1}^n \frac{1}{i}$$

- Theorem.** $H_n = \Theta(\log n)$
- Proof Idea. Upper and lower bound area under the curve



Card Guessing: Memoryfull

- Suppose we play the same game but now assume you have the ability to remember cards that have already been turned
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- Thus, $E[X] = \sum_{i=1}^n \frac{1}{n - i + 1} = \sum_{i=1}^n \frac{1}{i} = \Theta(\log n)$

Geometric Distribution

- Let's say we do a sequence of Bernoulli trials X_1, X_2, \dots where X_i where each trial is successful ($X_i = 1$) with probability p , and fails ($X_i = 0$) with probability $1 - p$
- **Question:** what is the expected number of trials until first success?
 - In expectation, what is the value of the first i such that $X_i = 1$?
 - E.g. number of coin flips until heads ($p = 1/2$)
 - E.g. number of times I roll a die until I get a 1 ($p = 1/6$)
- One way to solve it is to just do the sum:

$$\bullet \sum_{i=1}^{\infty} i(1-p)^{i-1}p$$



Geometric Expectation (using the sum)

$$\sum_{i=1}^{\infty} i(1-p)^{i-1}p = \sum_{i=1}^{\infty} \sum_{k=1}^i (1-p)^{i-1}p =$$

$$\sum_{k=1}^{\infty} \sum_{i=k}^{\infty} (1-p)^{i-1}p = \sum_{k=1}^{\infty} p(1-p)^{k-1} \sum_{i=0}^{\infty} (1-p)^i =$$

$$\sum_{k=1}^{\infty} p(1-p)^{k-1} \frac{1}{1-(1-p)} = \sum_{k=1}^{\infty} (1-p)^{k-1} = \sum_{k=0}^{\infty} (1-p)^k = \frac{1}{p}$$



Geometric Expectation (using the sum)

- Want to know, how many tries in expectation until first success
- Let's think about this recursively

$$X \leftarrow \begin{cases} 1 & \text{with prob. } p \\ 0 & \text{with prob. } (1 - p) \end{cases}$$

FindNumTries:

If $X = 1$

Return 1

If $X = 0$

Return $1 + \text{FindNumTries}$

If we fail in the first try, we start over from scratch!

- Let F be the number of times FindNumtries is called, what is $E(F)$?

Geometric Expectation (using the sum)

- Let F be the number of times FindNumtries is called, what is $E(F)$?
- $E(F) = E(F | X_1 = 1) \cdot \Pr(X_1 = 1) + E(F | X_1 = 0) \cdot \Pr(X_1 = 0)$
 $= (1 + 0) \cdot p + (1 + E(F)) \cdot (1 - p)$
- $E(F) = 1/p$

FindNumTries:

If $X = 1$

Return 1

If $X = 0$

Return $1 + \text{FindNumTries}$

If we fail in the first try, we start over from scratch!



Geometric Expectation: Formal Recursion

- Let X^p be a random variable indicating # flips until heads (with prob p)

- $$E(X^p) = \sum_{i=1}^{\infty} i(1-p)^{i-1}p$$

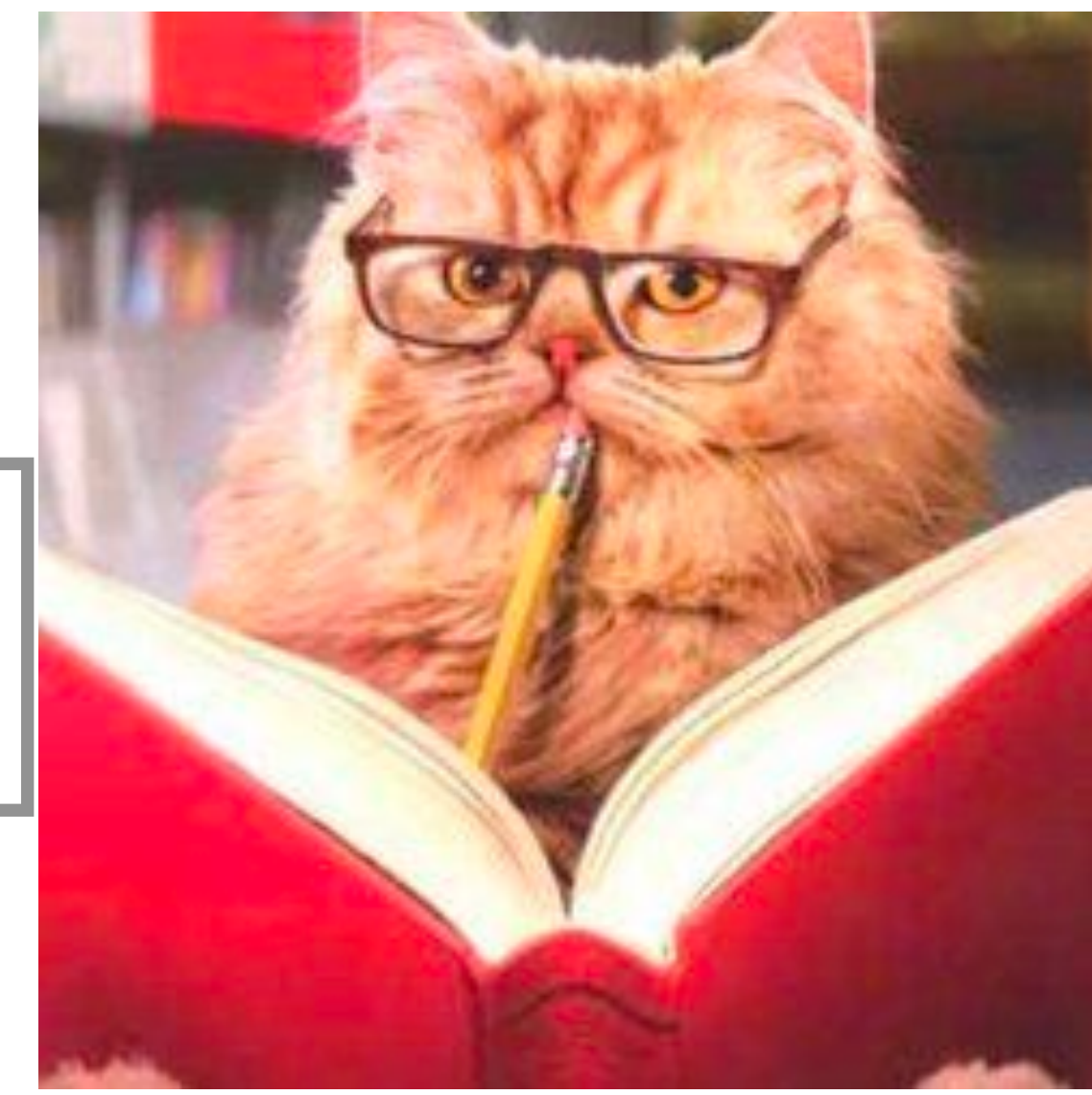
- We can then write

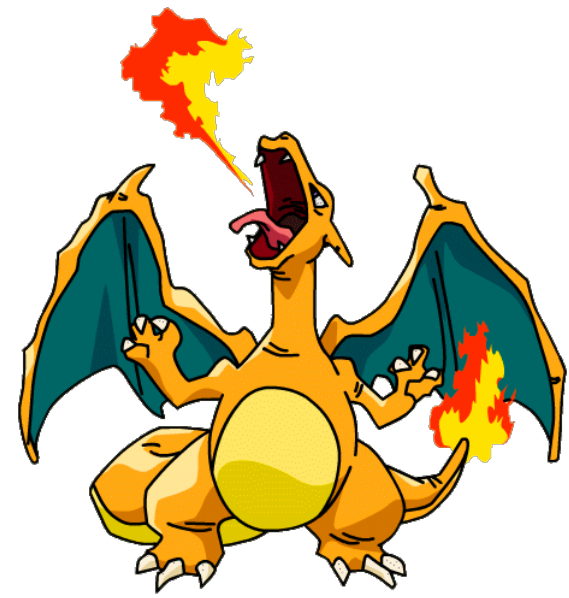
- $$E(X^p) = \sum_{i=1}^{\infty} i(1-p)^{i-1}p = p + \sum_{i=2}^{\infty} i(1-p)^{i-1}p = p + \sum_{i'=1}^{\infty} (1+i')(1-p)^{i'}p$$

- $$E(X^p) = p + (1-p) \sum_{i'=1}^{\infty} (1+i')(1-p)^{i'-1}p = p + (1-p)E(X^p + 1)$$

- $$E(X^p) = p + (1-p)(E(X^p) + 1)$$

You don't need to do this proof every time you use recursion. But, it can help if you're unsure of correctness



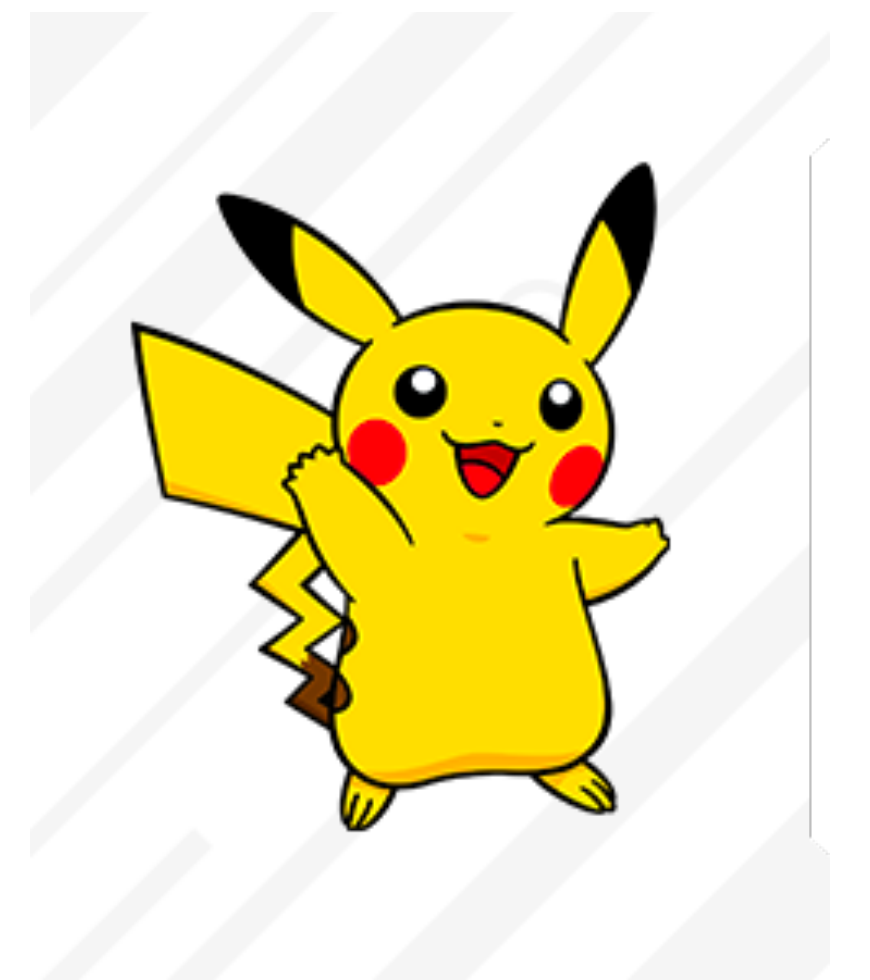


Coupon/Pokemon Collector Problem



Gotta' Catch 'Em All

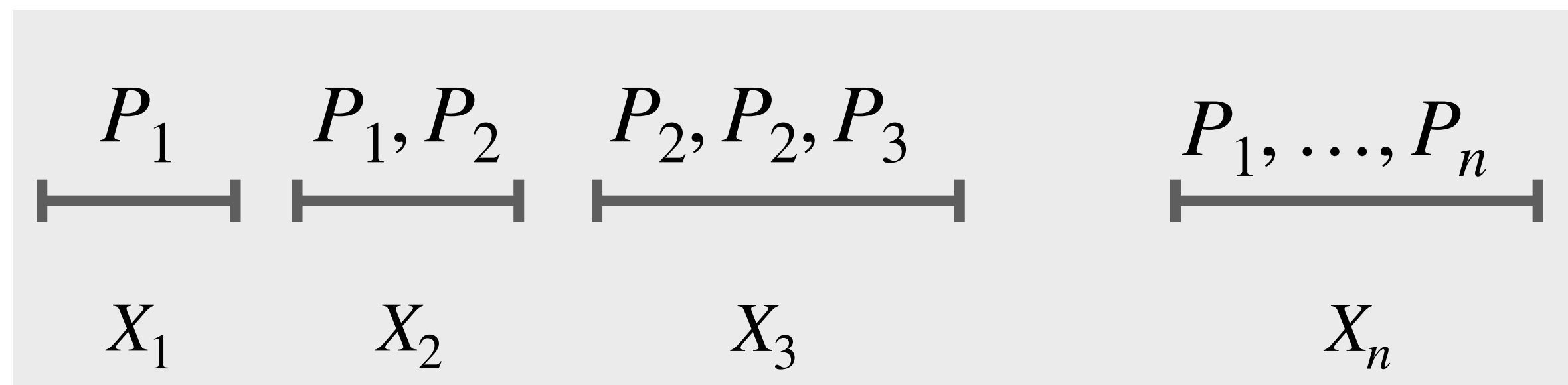
- Suppose there are n different types of Pokemon cards
- In each trial we purchase a pack that contains a Pokemon card
- We repeat until we have at least one of each type of card, how many packs does it take in expectation to collect all?
- Let X be the r.v. equal to the number of packs bought until you first have a card of each type. **Goal:** compute $E[X]$
- We break X into smaller random variables
- **Idea:** we make progress every time we get a card we don't already have



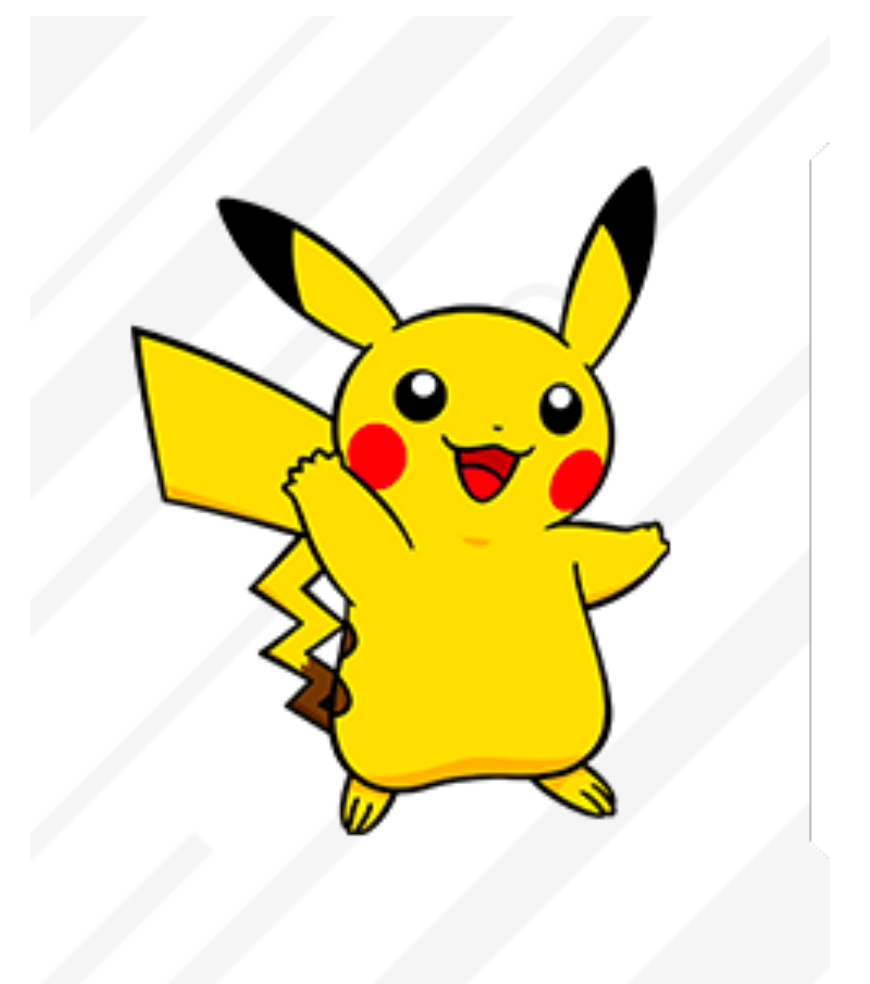
Pokemon Collector Problem

- Let X_i denote the "length of the i th phase", that is, the number of packs bought during the i th phase (i th phase ends as soon as we see the i th distinct card)

- Thus, $X = \sum_{i=1}^n X_i$

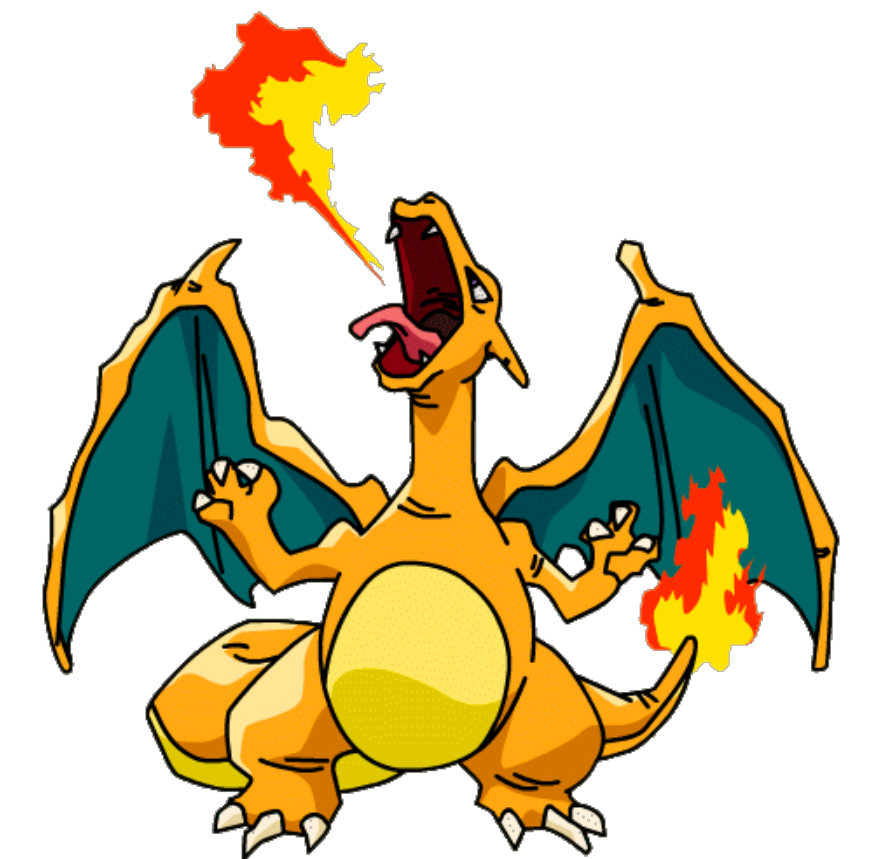
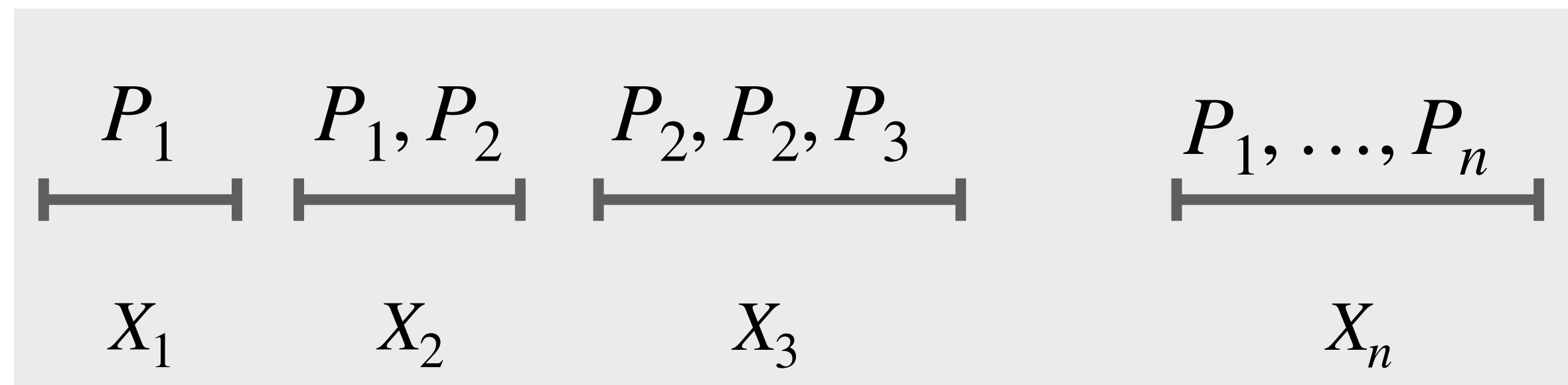


- Each phase can be thought of as flipping a biased coin until we see a head, where seeing a head = getting a new card



Pokemon Collector Problem

- $E[X_i]$ is the expected number of coin flips until success (expectation of a geometric r.v.)

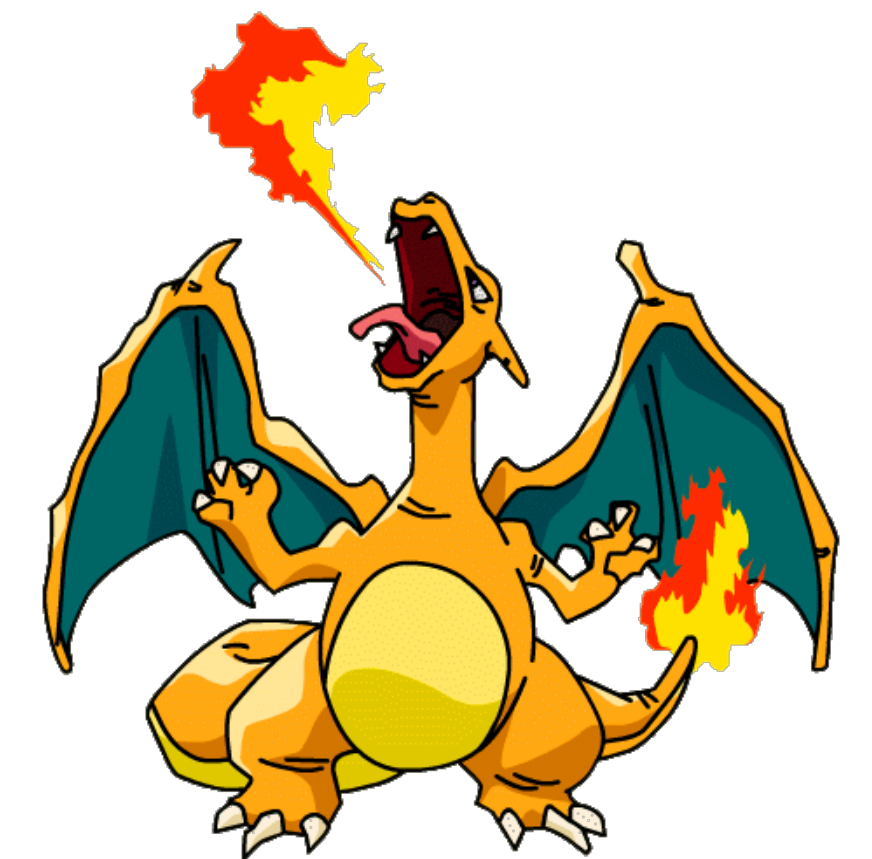
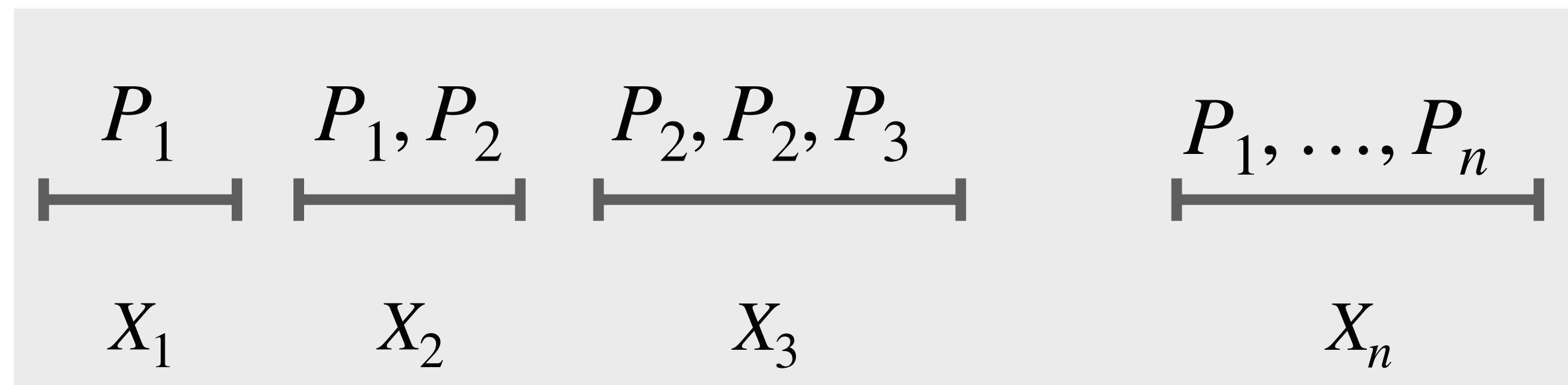


- We know, $E[X_i] = 1/p_i$ where p_i is the probability of success/ probability of seeing a heads during a coin flip in the i th phase
- Before the i th phase starts, we don't have $n - i + 1$ Pokemon
- Each of the n Pokemon are equally likely to be in a pack



Pokemon Collector Problem

- $E[X_i]$ is the expected number of coin flips until success
(expectation of a geometric r.v.)



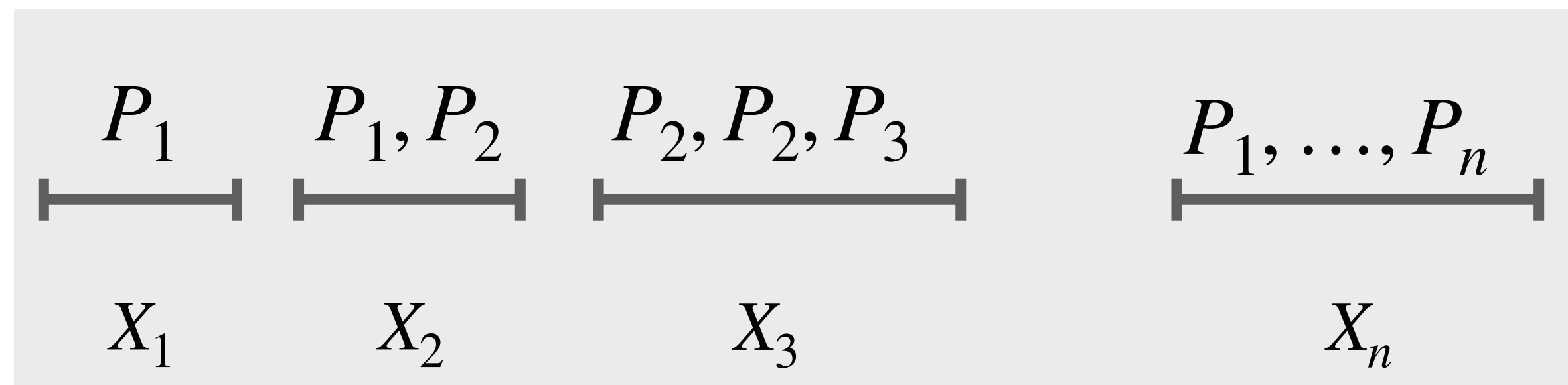
- We know, $E[X_i] = 1/p_i$ where p_i is the probability of success/
probability of seeing a heads during a coin flip in the i th phase

- $$p_i = \frac{n - i + 1}{n}$$



Pokemon Collector Problem

- We know, $E[X_i] = 1/p_i$ where p_i is the probability of success/
probability of seeing a heads during a coin flip in the i th phase



- $E[X_i] = \text{Expected}[\text{number of flips until first heads}] = 1/p_i = \frac{n - i + 1}{n}$
- $E[X] = E\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n E[X_i] = \sum_{i=1}^n \frac{n}{n - i + 1} = \sum_{i=1}^n \frac{n}{i} = nH_n = \Theta(n \log n)$

Acknowledgments

- Some of the material in these slides are taken from
 - Kleinberg Tardos Slides by Kevin Wayne (<https://www.cs.princeton.edu/~wayne/kleinberg-tardos/pdf/04GreedyAlgorithmsI.pdf>)
 - Jeff Erickson's Algorithms Book (<http://jeffe.cs.illinois.edu/teaching/algorithms/book/Algorithms-JeffE.pdf>)
 - Hamiltonian cycle reduction images from Michael Sipser's Theory of Computation Book