定向覆盖模糊测试工具的设计与实现

毕业设计中期检查

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What Fuzzing is?

Defination[1]

- **Fuzzing** Fuzzing is the execution of the PUT using input(s) sampled from an input space (the “fuzz input space”) that protrudes the expected input space of the PUT.
  - PUT: Program Under Test
- **Fuzz testing** Fuzz testing is the use of fuzzing to test if a PUT violates a correctness policy.
- **Fuzzer** A fuzzer is a program that performs fuzz testing on a PUT.
- **Bug Oracle** A bug oracle is a program, perhaps as part of a fuzzer, that determines whether a given execution of the PUT violates a specific correctness policy.
- **Fuzz Configuration** A fuzz configuration of a fuzz algorithm comprises the parameter value(s) that control(s) the fuzz algorithm.
- **Seed** A seed is a (commonly well-structured) input to the PUT, used to generate test cases by modifying it.

What Fuzzing is?

Fuzz Testing

Input: $\mathbb{C}$, t_{limit}
Output: $\mathbb{B}$ // a finite set of bugs

```
B ← ∅
C ← Preprocess(C)
while t_elapsed < t_limit ∧ Continue(C) do
    conf ← Schedule(C, t_elapsed, t_limit)
    tcs ← InputGen(conf)
    // O_bug is embedded in a fuzzer
    B′, execinfos ← InputEval(conf, tcs, O_bug)
    C ← ConfUpdate(C, conf, execinfos)
    B ← B ∪ B′
return B
```

Fuzzing Algorithm

Input: $\mathbb{C}$, t_{limit}
Output: $\mathbb{B}$ // a finite set of bugs
$\mathbb{B} \leftarrow \varnothing$
$\mathbb{C} \leftarrow$ Preprocess($\mathbb{C}$)
while $t_{elapsed} < t_{limit} \wedge$ Continue($\mathbb{C}$) **do**
    conf $\leftarrow$ Schedule($\mathbb{C}$, $t_{elapsed}$, t_{limit})
    tcs $\leftarrow$ InputGen(*conf*)
    // O_{bug} is embedded in a fuzzer
    $\mathbb{B}'$, execinfos $\leftarrow$ InputEval(*conf*, *tcs*, O_{bug})
    $\mathbb{C} \leftarrow$ ConfUpdate($\mathbb{C}$, *conf*, *execinfos*)
    $\mathbb{B} \leftarrow \mathbb{B} \cup \mathbb{B}'$
return $\mathbb{B}$

- $\mathbb{C}$:a set of fuzz configurations
- t_{limit}: timeout
- $\mathbb{B}$: a set of discovered bugs

Fuzzing Algorithm

Input: $\mathbb{C}$, t_{limit}
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```
𝔹 ← ∅
ℂ ← Preprocess(ℂ)
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```

$\text{PREPROCESS}(\mathbb{C}) \rightarrow \mathbb{C}$

- **Instrumentation**
  - grey-box and white-box fuzzers
  - static/dynamic(INPUTEVAL)
- **Seed Selection**
  - weed out potentially redundant configurations
- **Seed Trimming**
  - reduce the size of seeds
- **Preparing a Driver Application**
  - library Fuzzing, kernal Fuzzing

Fuzzing Algorithm

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return 𝔹
```

Stop Condition

- $t_{\text{elapsed}} < t_{\text{limit}}$
- CONTINUE $(\mathbb{C}) \rightarrow \{\texttt{True}, \texttt{False}\}$
  - Determine whether a new fuzz iteration should occur

Fuzzing Algorithm

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SCHEDULE ($\mathbb{C}$, $t_{elapsed}$, t_{limit}) $\rightarrow$ conf

- **Function**
  - Pick important information(conf)
- **FCS Problem**
  - *exploration*:Spent time on gathering more accurate information on each configuration to inform future decisions
  - *exploitation*:Spent time on fuzzing the configurations that are currently believed to lead to more favorable outcomes

Fuzzing Algorithm

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INPUTGEN (conf)$\longrightarrow$ tcs

- **function**
  - Generate testcases
- **classification**
  - Generation-based(*Model-based*)
  - Mutation-based(*Model-less*)
  - White-box Fuzzers: symbolic execution

Fuzzing Algorithm

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InputEval (conf, tcs, O_{bug})
$\rightarrow \mathbb{B}'$, execinfos

- **Fuzzing PUT**
  - tcs
  - $\mathbb{B}'$
- **Feedback Information**
  - conf, tcs
  - execinfos (tcs,crashes,stack backtrace hash,edge coverage,etc.)

Fuzzing Algorithm

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- ConfUpdate ($\mathbb{C}$, conf, execinfos) $\rightarrow \mathbb{C}$
  - Update Fuzz Configuration(distinguishablity)
  - Seed Pool Update
- $\mathbb{B} \cup \mathbb{B}' \rightarrow \mathbb{B}$
  - Update Bugs Set

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stop condition

- $t_{elapsed} < t_{limit}$
- Continue ($\mathbb{C}$) $\rightarrow$ {True, False}
  - Determine whether a new fuzz iteration should occur

Classification

The amount of collected information defines the color of a fuzzer[1].

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```

- **program instrumentation**
  - static
  - dynamic
- processor traces
- system call usage
- etc.

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```

Program Instrumentation

- Static
  - source code
  - intermediate code
  - binary-level
- Dynamic

Classification

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Program Instrumentation

- Static
- Dynamic
  - dynamically-linked libraries
  - execution feedback: branch coverage, new path, etc.
  - race condition bugs: thread scheduling

Classification

Classification of Fuzzing

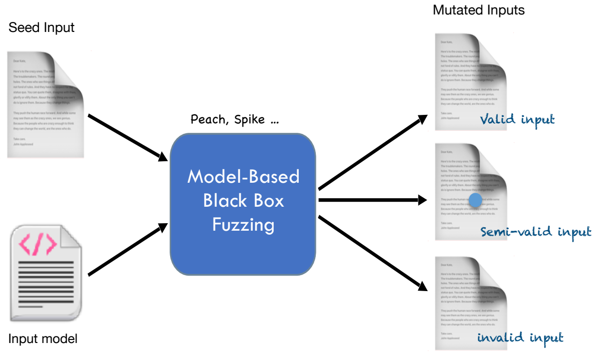
- **Black-box Fuzzing**
  - no program analysis, no feedback
- **White-box Fuzzing**
  - mostly program analysis
- **Grey-box Fuzzing**
  - no program analysis, but feedback

Why Grey-box Fuzzing ?

- Black-box Fuzzing

  Defination: techniques that do not see the internals of the PUT,and can observe only the input/output behavior of the PUT, treating it as a black-box[1].

  -No program analysis, no feedback



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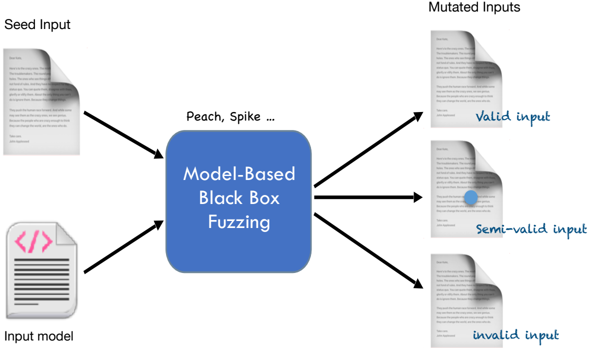


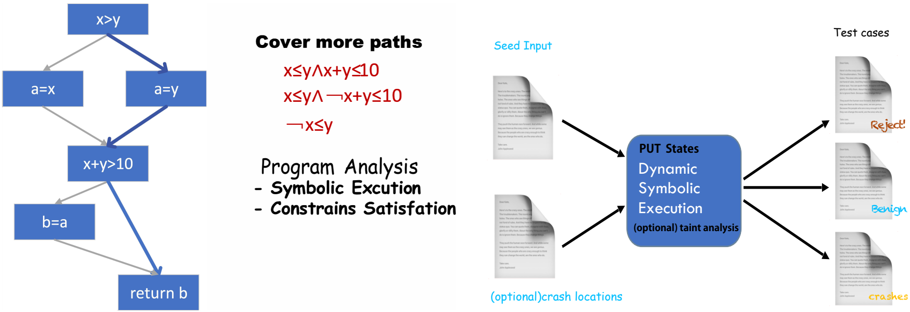
- You have no view of the PUT,but have some view of the input/output domain
- Fuzzing congfigurations are not changed according to some feedback
  - some fuzzer may add the testcases to seed pool
- Not effective

Why Grey-box Fuzzing ?

- White-box Fuzzing

  Defination: techniques that generates test cases by analyzing the internals of the PUT and the information gathered when executing the PUT[1].

  - Requires heavy-weight program analysis and constraint solving.



Why Grey-box Fuzzing ?

- White-box Fuzzing

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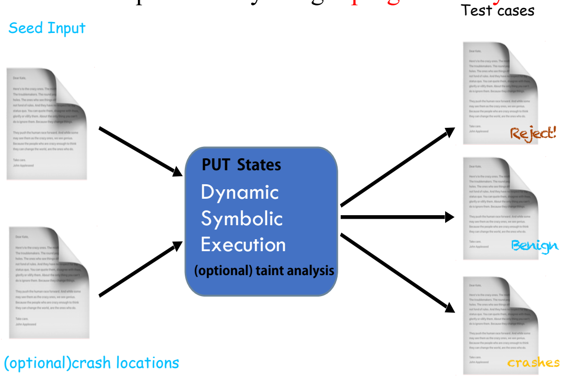
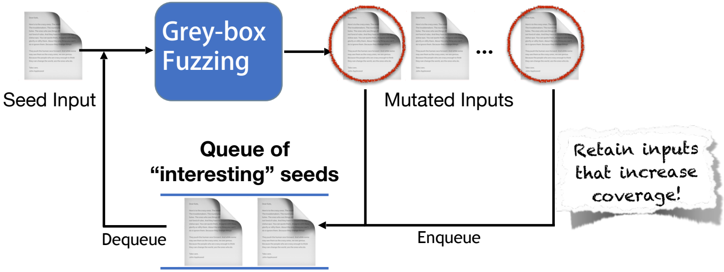


- You have the view of the PUT state(CFG,CG)
- Heavy-weight Program analysis (effective but not efficient!)

Why Grey-box Fuzzing ?

- Grey-box Fuzzing
  Defination: techniques that can obtain *some* information internal to the PUT and/or its executions to generates test cases[1].
  - Uses only lightweight instrumentation to glean some program structure
  - And coverage feedback



Why Grey-box Fuzzing ?

Grey-box Fuzzing is frequently used

- **State-of-the-art** in automated vulnerability detection
- **Extremely efficient** coverage-based input generation
  - All program analysis before/at instrumentation time.
  - Start with a seed corpus, choose a seed file, fuzz it.
  - Add to corpus only if new input increases coverage.

Why Directed Grey-box Fuzzing ?

Directed Fuzzing has many applications

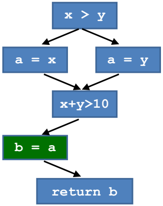
- **Patch Testing**: reach changed statements
- **Crash Reproduction**: exercise stack trace
- **SA Report Verification**: reach “dangerous” location
- **Information Flow Detection**: exercise source-sink pairs

Why Directed Grey-box Fuzzing ?

Directed Fuzzing

- **Goal:reach a specific target**
  - **Target Locations**: the line number in the source code or the virtual memory address at the binary level[2].
  - **Target Bugs**: use-after-free vulnerabilities, etc.

- **DSE:classical constraint satisfaction problem**
  - uses program analysis and constraint solving to generate inputs that systematically and effectively explore the state space of feasible paths[3].
  - **Program analysis** to identify **program paths** that reach given program locations.
  - **Symbolic Execution** to derive **path conditions** for any of the identified paths.
  - **Constraint Solving** to find an input

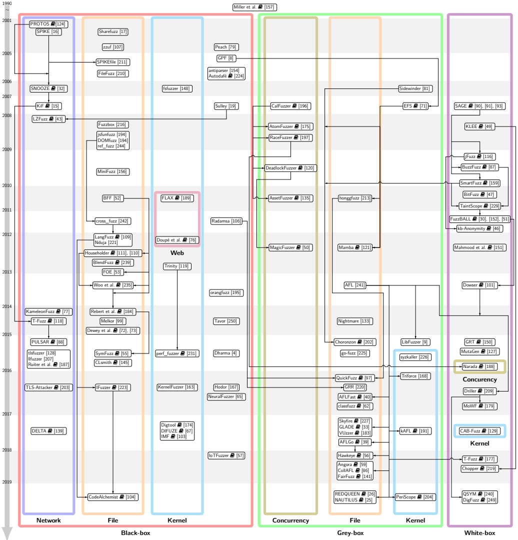


$$\varphi_1 = (x>y) \wedge (x+y>10)$$
$$\varphi_2 = \neg(x>y) \wedge (x+y>10)$$

Why Directed Grey-box Fuzzing ?

- **Effectiveness comes at the cost of efficiency**
- **Heavy-weight program analysis**

Genealogy tracing significant fuzzers' lineage[1]



[1]paper[1]-Figure1

Representative Work

- **AFLGo(2017)**[4]
- **Hawkeye(2018)**[5]
- **Github Repository**:awesome-directed-fuzzing
  `https://github.com/strongcourage/awesome-directed-fuzzing`

[1] MANÈS V J, HAN H, HAN C, et al. The art, science, and engineering of fuzzing: A survey[J]. IEEE Transactions on Software Engineering, 2019, 47(11) : 2312 – 2331.

[2] WANG P, ZHOU X, LU K, et al. The Progress, Challenges, and Perspectives of Directed Greybox Fuzzing[EB]. arXiv, 2022.

[3] MA K-K, YIT PHANG K, FOSTER J S, et al. Directed symbolic execution[C] // Static Analysis: 18th International Symposium, SAS 2011, Venice, Italy, September 14-16, 2011. Proceedings 18. 2011 : 95 – 111.

[4] BÖHME M, PHAM V-T, NGUYEN M-D, et al. Directed Greybox Fuzzing[C/OL] // Proceedings of the 2017 ACM SIGSAC Conference on Computer and Communications Security. Dallas Texas USA : ACM, 2017 : 2329 – 2344. http://dx.doi.org/10.1145/3133956.3134020.

[5] CHEN H, XUE Y, LI Y, et al. Hawkeye: Towards a Desired Directed Grey-Box Fuzzer[C/OL] // Proceedings of the 2018 ACM SIGSAC Conference on Computer and Communications Security. Toronto Canada : ACM, 2018 : 2095 – 2108. http://dx.doi.org/10.1145/3243734.3243849.

Thanks!