

1a. Prove that $(A \wedge B \rightarrow C \wedge D) \vdash (A \wedge B \rightarrow C)$ ("conjunctive rule splitting") is a **sound rule-of inference** using a truth table.

A rule R is sound iff for all sentences, a, b if $a \vdash b$ then $a \models b$

ANSWER

A	B	C	D	$A \wedge B$	$C \wedge D$	$A \wedge B \rightarrow C \wedge D$	$A \wedge B \rightarrow C$
T	T	T	T	T	T	T	T
T	T	T	F	T	F	F	T
T	T	F	T	T	F	F	F
T	T	F	F	T	F	F	F
T	F	T	T	F	T	T	T
T	F	T	F	F	F	T	T
T	F	F	T	F	F	T	T
T	F	F	F	F	F	T	T
F	T	T	T	F	T	T	T
F	T	T	F	F	F	T	T
F	T	F	T	F	F	T	T
F	T	F	F	F	F	T	T
F	F	T	T	F	T	T	T
F	F	T	F	F	F	T	T
F	F	F	T	F	F	T	T
F	F	F	F	F	F	T	T

We see that all models that satisfy $(A \wedge B \rightarrow C \wedge D)$ also satisfy $(A \wedge B \rightarrow C)$ and thus this is a sound rule of inference.

1b. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using Natural Deduction.

(Hint: it might help to use a ROI for "Implication Introduction". If you have a Horn clause, with 1 positive literal and n-1 negative literals, like $(\neg X \vee Z \vee \neg Y)$, you can transform it into a rule by collecting the negative literals as positive antecedents, e.g. $X \wedge Y \rightarrow Z$. This is a truthpreserving operation (hence sound), which you could prove to yourself using a truth table.)

ANSWER

1. $(A \wedge B \rightarrow C \wedge D)$ [Premise]
2. $C \wedge D$ [Premise]
3. $\neg(A \wedge B) \vee (C \wedge D)$ [Implication Elimination, 1]
4. $(\neg A \vee \neg B) \vee (C \wedge D)$ [De Morgan, 3]
5. $(\neg A \vee \neg B) \vee C$ [And Elimination, 4, 2]
6. $A \wedge B \rightarrow C$ [Implication Introduction, 5]

1c. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using Resolution.

We will use resolution refutation. First we convert our KB into CNF:

1. $(A \wedge B \rightarrow C \wedge D)$ [KB]
2. $\neg(A \wedge B) \vee (C \wedge D)$ [Implication Elimination, 1]
3. $(\neg A \vee \neg B) \vee (C \wedge D)$ [De Morgan, 2]
4. $(\neg A \vee \neg B \vee C) \wedge (\neg A \vee \neg B \vee D)$ [Distribution, 3]

Our new KB = $\{(\neg A \vee \neg B \vee C), (\neg A \vee \neg B \vee D)\}$

Now, we negate our query and add it to the KB

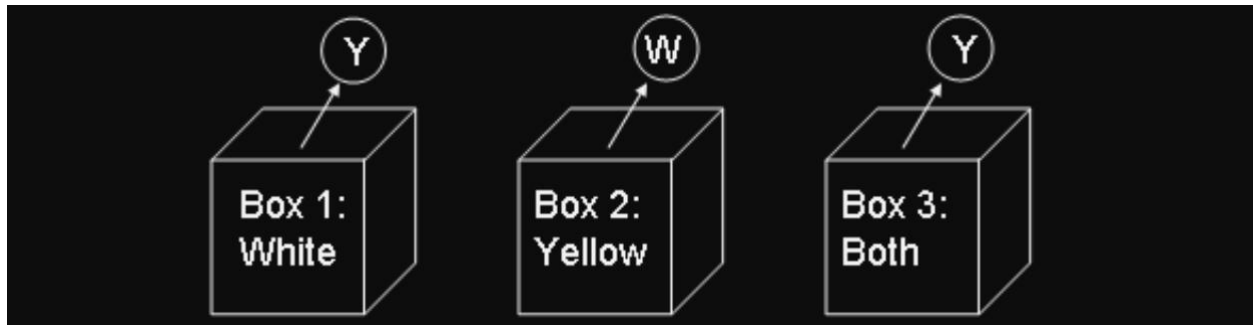
1. $\neg(A \wedge B \rightarrow C)$ [Negation]
2. $\neg(\neg(A \wedge B) \vee C)$ [Implication Elimination, 1]
3. $\neg\neg(A \wedge B) \wedge \neg C$ [DeMorgan, 2]
4. $A \wedge B \wedge \neg C$ [DN, 3]

Now our KB = $\{(\neg A \vee \neg B \vee C), (\neg A \vee \neg B \vee D), A, B, \neg C\}$. From this we should be able to derive an inconsistency.

1. A [KB]
2. B [KB]
3. $\neg C$ [KB]
4. $(\neg A \vee \neg B \vee C)$ [KB]
5. $(\neg B \vee C)$ [Resolve 1, 4]
6. (C) [Resolve 2, 5]
7. \emptyset [Resolve, 3, 6]

2. Sammy's Sport Shop

You are the proprietor of Sammy's Sport Shop. You have just received a shipment of three boxes filled with tennis balls. One box contains only yellow tennis balls, one box contains only white tennis balls, and one contains both yellow and white tennis balls. You would like to stock the tennis balls in appropriate places on your shelves. Unfortunately, the boxes have been labeled incorrectly; the manufacturer tells you that you have exactly one box of each, but that each box is definitely labeled wrong. You draw one ball from each box and observe its color. Given the initial (incorrect) labeling of the boxes above, and the three observations, use Propositional Logic to infer the correct contents of the middle box.



Use propositional symbols in the following form: O1Y means a yellow ball was drawn (observed) from box 1, L1W means box 1 was initially labeled white, C1W means box 1 contains (only) white balls, and C1B means box 1 actually contains both types of tennis balls. Note, there is no 'O1B', etc, because you can't directly "observe both". When you draw a tennis ball, it will either be white or yellow. The initial facts describing this particular situation are: {O1Y, L1W, O2W, L2Y, O3Y, L3B}

2a. Using these propositional symbols, write a propositional knowledge base (sammy.kb) that captures the knowledge in this domain (i.e. implications of what different observations or labels mean, as well as constraints inherent in this problem, such as that all boxes have different contents). Do it in a complete and general way, writing down all the rules and constraints, not just the ones needed to make the specific inference about the middle box. Do not include derived knowledge that depends on the particular labeling of this instance shown above; stick to what is stated in the problem description above. Your KB should be general enough to reason about any alternative scenario, not just the one given above (e.g. with different observations and labels and box contents).

ANSWER

KB = {

- a. $O1Y \rightarrow C1Y \vee C1B$
- b. $O1W \rightarrow C1W \vee C1B$
- c. $O2Y \rightarrow C2Y \vee C2B$
- d. $O2W \rightarrow C2W \vee C2B$
- e. $O3Y \rightarrow C3Y \vee C3B$
- f. $O3W \rightarrow C3W \vee C3B$
- g. $L1W \rightarrow (C1Y \vee C1B) \wedge \neg C1W$
- h. $L1Y \rightarrow (C1W \vee C1B) \wedge \neg C1Y$
- i. $L1B \rightarrow (C1Y \vee C1W) \wedge \neg C1B$
- j. $L2W \rightarrow (C2Y \vee C2B) \wedge \neg C2W$
- k. $L2Y \rightarrow (C2W \vee C2B) \wedge \neg C2Y$
- l. $L2B \rightarrow (C2Y \vee C2W) \wedge \neg C2B$
- m. $L3W \rightarrow (C3Y \vee C3B) \wedge \neg C3W$
- n. $L3Y \rightarrow (C3W \vee C3B) \wedge \neg C3Y$
- o. $L3B \rightarrow (C3Y \vee C3W) \wedge \neg C3B$
- p. $\neg(C1W \wedge C2W) \wedge \neg(C1W \wedge C3W) \wedge \neg(C2W \wedge C3W)$
- q. $\neg(C1Y \wedge C2Y) \wedge \neg(C1Y \wedge C3Y) \wedge \neg(C2Y \wedge C3Y)$
- r. $\neg(C1B \wedge C2B) \wedge \neg(C1B \wedge C3B) \wedge \neg(C2B \wedge C3B)$
- s. $\neg(L1W \wedge L2W) \wedge \neg(L1W \wedge L3W) \wedge \neg(L2W \wedge L3W)$
- t. $\neg(L1Y \wedge L2Y) \wedge \neg(L1Y \wedge L3Y) \wedge \neg(L2Y \wedge L3Y)$
- u. $\neg(L1B \wedge L2B) \wedge \neg(L1B \wedge L3B) \wedge \neg(L2B \wedge L3B)$
- v. $\neg(L1W \wedge L1Y) \wedge \neg(L1W \wedge L1B) \wedge \neg(L1Y \wedge L1B)$
- w. $\neg(L2W \wedge L2Y) \wedge \neg(L2W \wedge L2B) \wedge \neg(L2Y \wedge L2B)$
- x. $\neg(L3W \wedge L3Y) \wedge \neg(L3W \wedge L3B) \wedge \neg(L3Y \wedge L3B)$
- y. $\neg O1B \wedge \neg O2B \wedge \neg O3B$
- z. $\neg(O1W \wedge O1Y) \wedge \neg(O2W \wedge O2Y) \wedge \neg(O3W \wedge O3Y)$

}

a-f: Consequence of observables

g-o: Consequence of labels

p-r: Constraints (content of boxes)

s-x: Constraints (labels of boxes, one per box)

y-z: Constraints (observables)

2b. Prove that box 2 must contain white balls (C2W) using Natural Deduction.

Facts: {O1Y, L1W, O2W, L2Y, O3Y, L3B}

1. $O3Y \rightarrow C3Y \vee C3B$ [e]

2. $L3B \rightarrow (C3Y \vee C3W) \wedge \neg C3B$ [o]
3. $L3B$ [Fact]
4. $(C3Y \vee C3W) \wedge \neg C3B$ [MP, 2, 3]
5. $\neg C3B$ [AE, 2, 6]
6. $C3Y$ [Simp, 1, 5] (Established that box 3 is yellow)
7. $O1Y \rightarrow C1Y \vee C1B$
8. $O1Y$ [Fact]
9. $C1Y \vee C1B$ [MP, 7, 8]
10. $C1B$ [q, 6, 9] (Constraint) (Established that box 1 is both)
11. $O2W \rightarrow C2W \vee C2B$ [d]
12. $O2W$ [Fact]
13. $C2W \vee C2B$ [MP, 11, 12]
14. $C2W$ [r, 10, 13] (Constraint) (box 2 is white)

2c. Convert your KB to CNF.

CNF(KB) = {

1. $\neg O1Y \vee C1Y \vee C1B$
2. $\neg O1W \vee C1W \vee C1B$
3. $\neg O2Y \vee C2Y \vee C2B$
4. $\neg O2W \vee C2W \vee C2B$
5. $\neg O3Y \vee C3Y \vee C3B$
6. $\neg O3W \vee C3W \vee C3B$
7. $(\neg L1W \vee C1Y \vee C1B)$
8. $(\neg L1W \vee \neg C1W)$
9. $(\neg L1Y \vee C1W \vee C1B)$
10. $(\neg L1Y \vee \neg C1Y)$
11. $(\neg L1B \vee C1Y \vee C1W)$
12. $(\neg L1B \vee \neg C1B)$
13. $(\neg L2W \vee C2Y \vee C2B)$
14. $(\neg L2W \vee \neg C2W)$
15. $(\neg L2Y \vee C2W \vee C2B)$
16. $(\neg L2Y \vee \neg C2Y)$
17. $(\neg L2B \vee C2Y \vee C2W)$
18. $(\neg L2B \vee \neg C2B)$
19. $(\neg L3W \vee C3Y \vee C3B)$
20. $(\neg L3W \vee \neg C3W)$
21. $(\neg L3Y \vee C3W \vee C3B)$
22. $(\neg L3Y \vee \neg C3Y)$
23. $(\neg L3B \vee C3Y \vee C3W)$
24. $(\neg L3B \vee \neg C3B)$
25. $(\neg C1W \vee \neg C2W)$
26. $(\neg C1W \vee \neg C3W)$

- 27. $(\neg C2W \vee \neg C3W)$
- 28. $(\neg C1Y \vee \neg C2Y)$
- 29. $(\neg C1Y \vee \neg C3Y)$
- 30. $(\neg C2Y \vee \neg C3Y)$
- 31. $(\neg C1B \vee \neg C2B)$
- 32. $(\neg C1B \vee \neg C3B)$
- 33. $(\neg C2B \vee \neg C3B)$
- 34. $(\neg L1W \vee L2W)$
- 35. $\neg L1W \vee \neg L3W$
- 36. $\neg L2W \vee \neg L3W$
- 37. $\neg L1Y \vee \neg L2Y$
- 38. $\neg L1Y \vee \neg L3Y$
- 39. $\neg L2Y \vee \neg L3Y$
- 40. $\neg L1B \vee \neg L2B$
- 41. $\neg L1B \vee \neg L3B$
- 42. $\neg L2B \vee \neg L3B$
- 43. $\neg L1W \vee \neg L1Y$
- 44. $\neg L1W \vee \neg L1B$
- 45. $\neg L1Y \vee \neg L1B$
- 46. $\neg L2W \vee \neg L2Y$
- 47. $\neg L2W \vee \neg L2B$
- 48. $\neg L2Y \vee \neg L2B$
- 49. $\neg L3W \vee \neg L3Y$
- 50. $\neg L3W \vee \neg L3B$
- 51. $\neg L3Y \vee \neg L3B$
- 52. $\neg O1B$
- 53. $\neg O2B$
- 54. $\neg O3B$
- 55. $\neg O1W \vee \neg O1Y$
- 56. $\neg O2W \vee \neg O2Y$
- 57. $\neg O3W \vee \neg O3Y$

}

2d. Prove C2W using Resolution.

Facts: {O1Y, L1W, O2W, L2Y, O3Y, L3B} (also added to KB)

We add $\neg C2W$ to the KB.

1. $\neg O2W \vee C2W \vee C2B$ [KB, 4]
2. C2B [Resolve 1, facts]
3. $(\neg C1B \vee \neg C2B)$ [KB, 31]
4. $\neg C1B$ [Resolve 2,3]
5. $(\neg L1W \vee C1Y \vee C1B)$ [KB, 7]
6. C1Y [Resolve, facts, 4, 5]
7. $(\neg C1Y \vee \neg C3Y)$ [KB, 29]
8. $\neg C3Y$ [Resolve 6,7]
9. $(\neg L3B \vee C3Y \vee C3W)$ [KB, 23]
10. C3W [Resolve facts, 8, 9]
11. $(\neg C2W \vee \neg C3W)$ [KB, 27]
12. $\neg C2W$ [Resolve, 10, 11]
13. $(\neg L2Y \vee C2W \vee C2B)$ [KB, 15]
14. \emptyset [Resolve facts, 2, 12]

3. Do Forward Chaining for the CanGetToWork KB below. You don't need to follow the formal FC algorithm (with agenda/queue and counts array). Just indicate which rules are triggered (in any order), and keep going until all consequences are generated. Show the final list of all inferred propositions at the end. Is CanGetToWork among them?

KB = {

- a. $\text{CanBikeToWork} \rightarrow \text{CanGetToWork}$
- b. $\text{CanDriveToWork} \rightarrow \text{CanGetToWork}$
- c. $\text{CanWalkToWork} \rightarrow \text{CanGetToWork}$
- d. $\text{HaveBike} \wedge \text{WorkCloseToHome} \wedge \text{Sunny} \rightarrow \text{CanBikeToWork}$
- e. $\text{HaveMountainBike} \rightarrow \text{HaveBike}$
- f. $\text{HaveTenSpeed} \rightarrow \text{HaveBike}$
- g. $\text{OwnCar} \rightarrow \text{CanDriveToWork}$
- h. $\text{OwnCar} \rightarrow \text{MustGetAnnualInspection}$
- i. $\text{OwnCar} \rightarrow \text{MustHaveValidLicense}$
- j. $\text{CanRentCar} \rightarrow \text{CanDriveToWork}$
- k. $\text{HaveMoney} \wedge \text{CarRentalOpen} \rightarrow \text{CanRentCar}$
- l. $\text{HertzOpen} \rightarrow \text{CarRentalOpen}$
- m. $\text{AvisOpen} \rightarrow \text{CarRentalOpen}$
- n. $\text{EnterpriseOpen} \rightarrow \text{CarRentalOpen}$
- o. $\text{CarRentalOpen} \rightarrow \text{IsNotAHoliday}$
- p. $\text{HaveMoney} \wedge \text{TaxiAvailable} \rightarrow \text{CanDriveToWork}$
- q. $\text{Sunny} \wedge \text{WorkCloseToHome} \rightarrow \text{CanWalkToWork}$
- r. $\text{HaveUmbrella} \wedge \text{WorkCloseToHome} \rightarrow \text{CanWalkToWork}$

s. Sunny \rightarrow StreetsDry }

Facts: { Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity,
WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen }

ANSWER:

Irrelevant facts (negations are not allowed in FC): { Rainy, EnjoyPlayingSoccer, WorkForUniversity,
HertzClosed, McDonaldsOpen }

Relevant facts: {HaveMountainBike, WorkCloseToHome, HaveMoney, AvisOpen }

HaveMountainBike \rightarrow HaveBike

AvisOpen \rightarrow CarRentalOpen

CarRentalOpen \rightarrow IsNotAHoliday

HaveMoney \wedge CarRentalOpen \rightarrow CanRentCar

CanRentCar \rightarrow CanDriveToWork

CanDriveToWork \rightarrow CanGetToWork

Final list: {HaveBike, CarRentalOpen, IsNotAHoliday, CanRentCar, CanDriveToWork, CanGetToWork}

CanGetToWork is among them.

1. Do Backward Chaining for the CanGetToWork KB. In this case, you should follow the BC algorithm closely (the pseudocode for the propositional version of Back-chaining is given in the lecture slides). Important: when you pop a subgoal (proposition) from the goal stack, you should systematically go through all rules that can be used to prove it IN THE ORDER THEY APPEAR IN THE KB. In some cases, this will lead to back-tracking, which you should show. Also, the sequence of results depends on order in which antecedents are pushed onto the stack. If you have a rule like $A \wedge B \rightarrow C$, and you pop C off the stack, push the antecedents in reverse order, so B goes in first, then A; in the next iteration, A would be the next subgoal popped off the stack. 5. In what kinds of problems would it be better to use forward-chaining? When would it be better to use backward-chaining?

- CanGetToWork // initialize with query
- CanBikeToWork // pop CanGetToWork, push CanBikeToWork
- Sunny, WorkCloseToHome, HaveBike // pop CanBikeToWork, push antecedents in reverse order
- Sunny, WorkCloseToHome, HaveMountainBike // pop HaveBike, push antecedents
- Sunny, WorkCloseToHome // pop HaveMountainBike (fact)
- Sunny // pop WorkCloseToHome (fact)
- CanDriveToWork // (pop) Sunny is not provable, backtrack. Push CanDriveToWork
- OwnCar // pop CanDriveToWork, push antecedents
- CanRentCar // (pop) OwnCar is not provable, backtrack. Push CanRentCar
- CarRentalOpen, HaveMoney // pop CanRentCar, push antecedents in reverse order
- CarRentalOpen // pop HaveMoney (fact)
- HertzOpen // pop CarRentalOpen, push antecedents
- AvisOpen // (pop) HertzOpen is not provable, backtrack. Push AvisOpen
- \emptyset // pop AvisOpen (fact)

Forward-chaining is generally better when the goal or query is not clearly defined or there are multiple goals to be proven. Backward-chaining is generally better when there is a specific, well-defined goal or query to be proven.