

# MONADS FOR REAL PEOPLE

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# WHAT IS A MONAD?

- Monads define methods for chaining similar operations together and reducing boilerplate

# YOU'RE ALREADY USING MONADS

- Promise -> Javascript

- `$.get("/some/resources").then(function (val1) {  
 return $.get("/some/resources/" + val1[0]).then(function (val2) {  
 console.log("result of /some/resources/val1[0]", val2);  
 });  
});`
- `$.get("/some/resources").then(function (val1) {  
 return $.get("/some/resources/" + val1[0]);  
}).then(function (val2) {  
 console.log("result of /some/resources/val1[0]", val2);  
});`

# YOU'RE ALREADY USING MONADS

- Async -> C#

- ```
public Task<string> DoStuff()  
{  
    return Http.Get("/some/resources").ContinueWith(task =>  
    {  
        if (task.Failure) return task;  
        return Http.Get("/some/resources/" + task.Result);  
    });  
}
```
- ```
public async Task<string> DoStuff()  
{  
    var result = await Http.Get("/some/resources");  
    return await Http.Get("/some/resources/" + result);  
}
```

# YOU'RE ALREADY USING MONADS

- Try / ?. -> C#, Ruby

- C#:

- `if (foo != null && foo.Bar != null) {  
    return foo.Bar.Baz;  
}`
    - `foo?.Bar?.Baz`

- Ruby:

- `if foo.present? && foo.bar.present?  
    foo.bar.baz`
    - `foo.try(:bar).try(:baz) // try is from Rails`

# YOU'RE ALREADY USING MONADS

- List -> C#, Ruby, jQuery

- C#:

- `new [] { 1, 2, 3 }.ForEach(x => {  
 yield return x;  
 yield return x * 2;  
});`
    - `new [] { 1, 2, 3 }.SelectMany(x => new [] { x, x * 2 });`

- Ruby:

- `[1, 2, 3].map {|x| [x, x * 2]}.flatten`
    - `[1, 2, 3].flat_map {|x| [x, x * 2]}`

- jQuery: `$(".multiple-matches").find(".more-matches");`

# YOU'RE ALREADY USING MONADS

- The patterns we've just seen are all ad hoc or “one off” attempts to deal with types of computation chaining
- These types of problems are pretty common in programming, so it would be good to have a general solution
- *Monads* are one solution
- *Monads* provide a common pattern for handling computation chaining

# UPTOWN FUNCTOR

- A functor generally provides “context” around value types
  - Lists
  - Promises
  - Streams
  - Function objects (e.g. `Func<int,T>`)
  - Really any type which can support a generic parameter
- Technically, a functor must support *map*:
  - `F<a> -> (a -> b) -> F<b>`
  - e.g., `List<int> -> (int -> string) -> List<string>`
    - `new List<int>{ 1, 2, 3 }.Select(i => (i * 2).ToString())`



# CALL ME MAYBE

- **Maybe:** a way of representing a possible value or the absence of a value
  - Similar to `Nullable<>` generic type in C#
  - in Haskell, **Just 5** or **Nothing**
  - in F#, **Some 5** or **None** (*Maybe* is called the *Option* type in F#)
  - a better way of dealing with *null*

# CALL ME MAYBE

How many times have you wanted to write this:

- ```
var x = mightReturnNull(1);  
var y = mightAlsoReturnNull(x);  
var z = couldBeNullToo(y);  
return finalComputationOn(x, y, z);
```

But to be safe you have to write this:

- ```
var x = mightReturnNull(1);  
if (x == null) return null;  
var y = mightAlsoReturnNull(x);  
if (y == null) return null;  
var z = couldBeNullToo(y);  
if (z == null) return null;  
return finalComputationOn(x, y, z);
```

# CALL ME MAYBE

Can we have all the safety of

- ```
var x = mightReturnNull(1);  
if (x == null) return null;  
var y = mightAlsoReturnNull(x);  
if (y == null) return null;  
var z = couldBeNullToo(y);  
if (z == null) return null;  
return finalComputationOn(x, y, z);
```

without the boilerplate?

Yes!

- ```
maybe {  
  let! x = mightReturnNone 1  
  let! y = mightAlsoReturnNone x  
  let! z = couldBeNoneToo y  
  return finalComputationOn x y z  
}
```
- This looks like the original example, but has all the safety of the second - without the clutter
- Many monads are great at drawing out the intent of code and suppressing the necessary “plumbing”

# THE CHOICE OF A NEW GENERATION

- Choice:
  - One of two (or more) values, for example a function might return a data structure or an error message
  - called Choice in F#, Either in Haskell
  - you can think of it like *Option*, but with an error value instead of *None*
  - For our purposes:
    - ```
public struct Choice<TFail, TSuccess>
{
    public bool IsSuccess { get; }
    public TSuccess SuccessValue { get; }
    public TFail FailValue { get ; }
}
```

# THE CHOICE OF A NEW GENERATION

- ```
public Choice<string, SuccessStructure> DoSomeStuff() {  
    Choice<string, int> x = mightFail(1);  
    if (!x.IsSuccess) return Choice.Failure<string, SuccessStructure>(x.FailValue);  
    Choice<string, bool> y = mightAlsoFail(x);  
    if (!y.IsSuccess) return Choice.Failure<string, SuccessStructure>(y.FailValue);  
    Choice<string, string> z = couldBeFailureToo(y);  
    if (!z.IsSuccess) return Choice.Failure<string, SuccessStructure>(z.FailValue);  
    return finalComputationOn(x.SuccessValue, y.SuccessValue, z.SuccessValue);  
}
```
- ```
let doSomeStuff() =  
    choice {  
        let! x = mightFail 1  
        let! y = mightAlsoFail x  
        let! z = couldBeFailureToo y  
        let! ret = finalComputationOn x y z  
        return ret  
    }
```

# OKAY, BUT WHAT'S A MONAD?

- A Functor can be a monad by supporting two functions
  - `bind`
  - `return`
- **return** - take a raw value and put it in a Functor “context”
- **a -> M<a>**
  - `return 5 -> Some 5` // Option (Maybe)
  - `return “hi” -> Choice20f2 “hi”` // Choice (Either)
  - `return true -> [true]` // List

# MONAD IS AS MONAD DOES

- **bind**: connects one operation to the next
  - very similar to functor map, but where the transform can produce its own context
- **bind**:  $M\langle a \rangle \rightarrow (a \rightarrow M\langle b \rangle) \rightarrow M\langle b \rangle$ 
  - all functors involved must be the same type of functor (e.g. List, Option, Choice)
- e.g.  $\text{Option}\langle \text{int} \rangle \rightarrow (\text{int} \rightarrow \text{Option}\langle \text{float} \rangle) \rightarrow \text{Option}\langle \text{float} \rangle$
- shorthand operator is  $\gg=$  in Haskell
- So how do we define **bind** for *Option*?
  - `let optionBind (opt:Option<'a>) (func:('a -> Option<'b>)) : Option<'b> = ???`
    - ' means generic type in F#

# MONAD IS AS MONAD DOES

- ```
let optionBind opt funcValToOpt =  
    if opt.IsNone then  
        None  
    else  
        funcValToOpt opt.Value
```
- ```
let printIfOdd x =  
    if x % 2 = 1 then Some (x.ToString()) else None
```
- ```
optionBind (Some 3) printIfOdd -> Some "3"
```
- ```
optionBind (Some 8) printIfOdd -> None
```
- ```
optionBind None printIfOdd -> None
```



# MONAD IS AS MONAD DOES

- Let's return to our Option example:
  - `maybe {`  
    `let! x = mightReturnNone 1`  
    `let! y = mightAlsoReturnNone x`  
    `let! z = couldBeNoneToo y`  
    `return finalComputationOn x y z`  
    `}`
- What does the “unrolled” version look like?
  - `optionBind mightReturnNone(1) (fun x ->`  
    `optionBind mightAlsoReturnNone(x) (fun y ->`  
        `optionBind couldBeNoneToo(y) (fun z ->`  
            `optionReturn (finalComputationOn x y z)`  
        `)))`

# I FOUGHT THE LAW(S) AND THE LAW(S) WON

- Identity laws establish a symmetrical relationship between **bind** and **return**
- Associativity law guarantees predictable composition of operations

# I FOUGHT THE LAW(S) AND THE LAW(S) WON

## Left Identity

- `return x bind f` equals `f x`
- given:
  - `let printIfOdd x =  
    if (x % 2) == 1 then Some (x.ToString()) else None`
- `return 5 (i.e. Some 5) bind printIfOdd` equals `printIfOdd 5`
- as a counter example, if we
  - `let return v = Some (v+1) // for numeric types`
  - then `return 5 (i.e Some 6) bind printIfOdd` would equal `printIfOdd 6`

# I FOUGHT THE LAW(S) AND THE LAW(S) WON

## Right Identity

- `m bind return` equals `m`
- `Some 5 bind return` equals `Some 5`
- again if we
  - `let return v = Some (v+1) // for numeric types`
- `Some 5 bind return` would equal `Some 6`

# I FOUGHT THE LAW(S) AND THE LAW(S) WON

## Associativity – monadic composition

- say we have `a -> m<b>` and `b -> m<c>` which together produce `a -> m<c>`
  - e.g., `int -> Option<string>` and `string -> Option<boolean>` combines to form `int -> Option<boolean>`
    - ```
let f1 i : string =  
    if i > 5 then Some (i.ToString()) else None  
let f2 s : bool =  
    if s.Length > 2 then Some true else None  
let f1And2 = (fun a -> (f1 a) |> Option.bind f2)  
f1And2 is written as f1 >=> f2
```

# I FOUGHT THE LAW(S) AND THE LAW(S) WON

## Associativity

- to satisfy associativity,  
 $(aToMb \gg bToMc) \gg cToMd$  is equal to  
 $aToMb \gg (bToMc \Rightarrow cToMd)$
- $\gg$  - Kleisli composition of monadic functions
  - *not to be confused with Khaleesi composition*



# THE STATE THAT I AM IN

- The state monad allows the consumption / mutation of state to produce a series of outputs and a final state
- Example: given random number generation *randGen*
  - *maxValue -> seed -> randomValue \* newSeed*
    - *A \* B* means a tuple of A and B
    - because our data is immutable, we will output a new seed value with each invocation

# THE STATE THAT I AM IN

- Let's say we need three random values with different bounds

- `let rand1, seed1 = randGen 10 initialSeed`  
`let rand2, seed2 = randGen 25 seed1`  
`let rand3, _ = randGen 400 seed2 // we don't need the final seed`  
`[ rand1 ; rand2 ; rand3 ]`

- Using the state monad:

```
state {  
  let! rand1 = randGen 10  
  let! rand2 = randGen 25  
  let! rand3 = randGen 400  
  return [rand1 ; rand2 ; rand3]  
} initialSeed
```



# THE STATE THAT I AM IN

- Let's take a closer look at what is happening here:

```
state {  
  let! rand1 = randGen 10  
  let! rand2 = randGen 25  
  let! rand3 = randGen 400  
  return [rand1 ; rand2 ; rand3]  
} initialSeed
```

- You'll notice that if randGen is
  - `let randGen (maxValue: int) -> (seed:int) -> double * int`
  - we did not provide the final argument in the monad expression
  - bound expressions always produce the monadic type

# THE STATE THAT I AM IN

- The state monad is a function!
  - **State<T> = state -> T\*state**
  - state builds a pipeline somewhat similar to Kleisli composition
  - state value is replaced continuously as it threads through pipeline
  - `let return v = (fun s -> v*s)`
  - `let stateBind m:(`s -> `a*`s) f:(`a -> (`s -> `b*`s)) : (`s -> `b*`s) =  
 (fun origState ->  
 let value, newState = m origState  
 f value newState  
 )`
  - `// randGen(10) will receive initialSeed  
stateBind randGen(10) (fun rand1 ->  
 // randGen(25) will receive seed1 from randGen 10 initialSeed  
 stateBind randGen(25) (fun rand2 ->  
 // randGen(400) will receive seed2 from randGen 25 seed1  
 stateBind randGen(400) (fun rand3 ->  
 stateReturn [rand1 ; rand2 ; rand3]  
 )) initialSeed`

# THE STATE THAT I AM IN

- Using partial application we can produce a `State<T>` function from a normal function
  - `let stateBind m:(s -> a*s) f:(a -> (s -> b*s)) : (s -> b*s)`
  - note that you can normally define your state functions `a -> (s -> b*s)` as `a -> s -> b*s` (i.e. two arguments and a tuple for an output)
  - that is, this declaration:
    - `let doStuff x = (fun y -> ...)`
  - *equals* this declaration:
    - `let doStuff x y = ...`

# THE MORE YOU KNOW (READER)

- What if we need to have access to “global” data that doesn’t change?
- **env** → **T**
- lookup: name → phonebook → number
- reader {
  - let! num1 = lookup “Joe”
  - let! num2 = lookup “Jenny”
  - let! num3 = lookup “Phil”
  - return [num1 ; num2 ; num3]} phonebook
- let return v = (fun env → v)
- let readerBind m:(‘env → ‘a) f:(‘a → (‘env → ‘b)) : (‘env → ‘b) =
  - (fun env →
    - let aVal = m env
    - f aVal)

# LISTS (OR, THE WORLD IS FLAT)

- List:

- **bind:** `List<T> -> (T -> List<U>) -> List<U>`
- `flat_map / SelectMany`
- can be chained repeatedly:
- `list {`

```
    let! eachElem1 = generateAList
```

```
    let! eachElem2 = aListForEach eachElem
```

```
    let! eachElem3 = anotherListForEach eachElem2
```

```
    // if not for monads, we would have a List<List<List<x>>>
```

```
    return (eachElem3 * 10)
```

```
}
```

# ASYNCHRONICITY

- Async:

- **bind:** `Async<T> -> (T -> Async<U>) -> Async<U>`
- C# 4.5's *async*
- Javascript Promise chaining
- ```
let downloadAndExtractLinks url =  
    async {  
        let webClient = new System.Net.WebClient()  
        let html = webClient.DownloadString(url : string)  
        let! links = extractLinksAsync html  
        return url, links.Count  
    }
```

# MIXING IT UP: DESERIALIZATION

- Fundamentally, deserialization is
  - sequence of bytes -> DataStructure
  - for these examples, let's assume DataStructure =

```
{  
  Foo : int  
  Bar : string  
  Baz : Point  
}
```
- consuming the bytes as we read them sounds like a job for “state”

# MIXING IT UP: DESERIALIZATION

- `assume unpack<T>: byte[] -> Tuple(T * byte[])`
  - `state {`
    - `let! fooValue = unpack`
    - `let! barValue = unpack`
    - `let! bazValue = unpack`
    - `return {`
      - `Foo = fooValue`
      - `Bar = barValue`
      - `Baz = bazValue`
    - `}`
  - `} initialBytes`
- notice the high degree of type inference



# MIXING IT UP: DESERIALIZATION

- What if something goes wrong?
  - assume `unpack<T>: byte[] -> Choice<Tuple(T * byte[]), Error>`
  - **protectedState** {
    - `let! fooValue = unpack`
    - `let! barValue = unpack`
    - `let! bazValue = unpack`
    - `return {`
      - `Foo = fooValue`
      - `Bar = barValue`
      - `Baz = bazValue``}`
  - `} initialBytes`

# MIXING IT UP: DESERIALIZATION

- What if unpack requires a lookup table for deserializing subtypes?
  - assume *unpack*<T>: table -> byte[] -> Choice<Tuple(T \* byte[]),Error>
  - **readerProtectedState** {
    - let! fooValue = unpack
    - let! barValue = unpack
    - let! bazValue = unpack
    - return {
      - Foo = fooValue
      - Bar = barValue
      - Baz = bazValue
- } table initialBytes

(safe) Magic!



# REAL WORLD EXAMPLE 1: MULTI-LAYERED FUNCTOR CONTEXT

- Imagine that a program maintains many simultaneous communication channels, and logs each channel into its own file
- Each log file contains several lines
- Each line contains a timestamp and a JSON payload
- *grep*'ing to find interesting data works, but it has limitations
- merely extracting matching substrings can make it easy to lose context

# REAL WORLD EXAMPLE 1: MULTI-LAYERED FUNCTOR CONTEXT

- instead, use a functor `Log<data>`
- transforms that do functor mapping will maintain filename and timestamps
  - **filter:** keep all lines where payload matches filter
  - **timespan:** keep all times in time window
  - **select:** JSON path transform of payload
  - **merge:** combine two Logs
  - **mergeSingleton:** intersperse a Log file into other Log files
- also support writing out altered Log to a directory
- <https://github.com/willryan/TimestampedJsonLogFilter>

## REAL WORLD EXAMPLE 2: UNHAPPY PATHS

- Imagine going through a long, asynchronous process of building up a complicated data structure
- failures could happen at multiple points
- in C#, use a `Task<Choice<success,error>>` return type
- C# helps you with async syntactic sugar, but you're on your own for Choice type
  - I did add a “BindFail” for my choice type, for the rare case where you want to do something interesting on the fail case
    - technically not monadic, but oh well

# REAL WORLD EXAMPLE 3: STATEFUL REBOOT

- For another case, imagine code which must reboot multiple devices in a complicated order
  - rebooting is asynchronous
  - order is complex and next steps require analysis of current state
  - sounds like a job for async + state!
- State object has fluent methods for transform
  - indicate device was rebooted
    - `var state3 = state2.TriedReboot(new [] { device });`
  - add devices expected to be present after reboot
  - add which devices have been re-discovered after a reboot
    - `var state4 = state3.AddReady(newDevices);`
  - track number of tries to reboot a particular device
- I called `Bind()` “`ThenWith()`” to simplify nomenclature
- Still looks like callback hell in C# though

# BONUS: CSHARPMONAD!

- Use *LINQ* syntatic sugar to achieve “pretty” monads in C#
- Remember this?

- ```
public Choice<SuccessStructure, string> DoSomeStuff() {  
    int x = mightFail(1);  
    if (x.Fail) return Choice.Failure<SuccessStructure,string>(x.FailValue);  
    bool y = mightAlsoFail(x);  
    if (y.Fail) return Choice.Failure<SuccessStructure,string>(y.FailValue);  
    string z = couldBeFailureToo(y);  
    if (z.Fail) return Choice.Failure<SuccessStructure,string>(z.FailValue);  
    return finalComputationOn(x,y,z);  
}
```

- How about:

- ```
public Either<Output,string> doSomeStuff()  
{  
    from x in mightFail(1)  
    from y in mightAlsoFail(x)  
    from z in couldBeFailureToo(y)  
    select finalComputation(x,y,z);  
}
```



# BONUS: CSHARPMONAD!

- Because partial application is not available, state functions must be declared like
  - `public Func<state,Tuple<U,state>> stateFunc(T,OtherInputs)`
  - but you can write  
`public Tuple<U,state> stateFunc(T t, OtherInputs o, state s) and  
Func<state,Tuple<U,state>> stateFuncPrime(T t, OtherInputs o) {  
 return (st => stateFunc(t, o, st));  
}`
  - or use a technique such as in [MethodToDelegate](#) :)

# RESOURCES / FURTHER READING

<https://github.com/willryan/MonadsForRealPeople>

<http://learnyouahaskell.com/>

<http://fsharpforfunandprofit.com/posts/monadster/>

```
main = do
```

```
    questions <- getLine
```

```
    putStrLn "Thanks!"
```