

Date: 2015/12/01

Class:

ID:

Name:

- (10%) Please briefly describe the **Merge Sort** algorithm, and analyze its worst-case time complexity for sorting n numbers with the **Recursion Tree** method.
- (10%) Solve each following **recurrence** for $T(n)$.
 - $T(n) = 4T(n/2) + n$ (3%)
 - $T(n) = 4T(n/2) + n^2$ (3%)
 - $T(n) = 4T(n/2) + n^3$ (4%)
- (10%) $1234^{2345} \bmod 100 = ?$ Calculate by hand (without any calculator) efficiently.
- (15%) Consider the following recursive function where the global variable *count* is initialized to 0 and input n is a positive integer.

- Find out the asymptotic time complexity in Θ -notation. (5%)
- Find out the asymptotic space complexity in Θ -notation. (5%)
- Express the final value of *count* as a function of n . (5%)

```

Rec-x( n )
{
    if ( n = 1 ) or ( n = 2 ) then
        count ← count + 1
    else
    {
        Rec-x( n - 2 )
        Rec-x( n - 2 )
        Rec-x( n - 1 )
        count ← count + 1
    }
}

```

(10%) A **linear** consecutive- k -out-of- n system consists of n nodes arranged in a line, where the system fails if and only if some k consecutive nodes fail. A **circular** consecutive- k -out-of- n system consists of n nodes arranged on a circle, where the system fails if and only if some k consecutive nodes fail. Suppose the nodes are statistically independent and the reliability of node i is p_i for any $i \in \{1, 2, \dots, n\}$ (node i functions with probability p_i and fails with probability $1 - p_i$). Let $R_L(i, j, k)$ denote the reliability of the linear consecutive- k -out-of- n subsystem consisting of nodes $i, i + 1, i + 2, \dots, j$. $R_C(i, j, k)$ denote the reliability of the circular consecutive- k -out-of- n subsystem consisting of nodes $i, i + 1, i + 2, \dots, j$.

- Express $R_L(1, n, k)$ with $R_L(1, n-1, k)$ and $R_L(1, n-k-1, k)$. (5%)
- Express $R_C(1, n, k)$ with some $R_L(1, ?, k)$ terms. (5%)

$$\begin{array}{r}
 56 \\
 36 \\
 \hline
 36 \\
 8 \\
 \hline
 16 \\
 96 \\
 34 \\
 \hline
 96 \\
 4 \\
 \hline
 4 \\
 36 \\
 56 \\
 16 \\
 \hline
 0 \\
 16 \\
 \hline
 16 \\
 345 \\
 64 \\
 8 \\
 \hline
 44
 \end{array}$$

$$\begin{array}{r}
 56 \\
 56 \\
 36 \\
 \hline
 2 \\
 36 \\
 36 \\
 16 \\
 \hline
 8 \\
 96
 \end{array}$$

$$\begin{array}{r}
 96 \\
 96 \\
 \hline
 76 \\
 4 \\
 \hline
 16
 \end{array}$$

$$\begin{array}{r}
 2345 \\
 1044 \\
 \hline
 0 \\
 2345 \\
 1048 \\
 \hline
 2917 \\
 256 \\
 \hline
 41 \\
 32 \\
 \hline
 9
 \end{array}$$

6. (10%) What are the main differences between **Randomized Algorithms** and Non-Randomized Algorithms?
7. (15%) Let m and n be positive integers such that $1 < m < n$. Given a procedure `Unbiased_RBG` that outputs an unbiased random bit which will be 0 or 1 with equal probability, $1/2$, please
- Design an algorithm, `Biased_RBG`, with calls of subroutine `Unbiased-RBG` to output 1 with probability $m/(m+n)$ and 0 with probability $n/(m+n)$. (5%)
 - Prove that your algorithm is correct. (5%)
 - What's the expected running time of your algorithm as a function of m and n ? (5%)
8. (10%) Consider the following program segment.

```

i ← 2
x ← 3
while ( i < n ) do
{
    i ← i × i
    x ← x + 1
}

```

- Find the time complexity (in Θ notation) of this program segment. (5%)
 - Express the final value of x as a function of n . (5%)
9. (10%) Write the `SELECT()` algorithm to output the i -th smallest element of the input array in worst-case time complexity $O(n)$ when given an array of n elements and an index integer i .

$$\begin{array}{r} 34 \\ 34 \\ \hline 36 \\ 2 \\ \hline 56 \end{array}$$

$$\begin{array}{r} 34 \\ 34 \\ \hline 36 \\ 2 \\ \hline 56 \end{array}$$

$$\begin{array}{r} 96 \\ 16 \\ \hline 76 \\ 6 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 96 \\ 16 \\ \hline 76 \\ 6 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 56 \\ 56 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 56 \\ 56 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 36 \\ 36 \\ \hline 16 \\ 8 \\ \hline 96 \end{array}$$

$$\begin{array}{r} 76 \\ 76 \\ \hline 16 \\ 8 \\ \hline 96 \end{array}$$

$$\begin{array}{r} 36 \\ 56 \\ \hline 16 \\ 0 \\ \hline \end{array}$$

$$\begin{array}{r} 36 \\ 56 \\ \hline 16 \\ 0 \\ \hline \end{array}$$

$$\begin{array}{r} 96 \\ 96 \\ \hline 76 \\ 4 \\ \hline 16 \end{array}$$

$$\begin{array}{r} 96 \\ 96 \\ \hline 76 \\ 4 \\ \hline 16 \end{array}$$

$$\begin{array}{r} 16 \\ 96 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 36 \\ 34 \\ \hline 44 \\ 8 \\ \hline 24 \end{array}$$

$$\begin{array}{r} 16 \\ 96 \\ \hline 96 \\ 4 \\ \hline 36 \end{array}$$

$$\begin{array}{r} 16 \\ 16 \\ \hline 96 \\ 6 \\ \hline 56 \end{array}$$

$$\begin{array}{r} 16 \\ 16 \\ \hline 96 \\ 6 \\ \hline 56 \end{array}$$

$$\begin{array}{r} 36 \\ 34 \\ \hline 44 \\ 8 \\ \hline 24 \end{array}$$