

(上列各項，考生須逐項填明，違者該科試卷以零分計)

Please fill out the above information, otherwise the score of this quiz will be zero.

1. $f(n) = O(n^{\log_b a - \epsilon}) \quad \epsilon > 0$
 $\rightarrow O(n^{\log_b a})$

$f(n) = \Theta(n^{\log_b a} \lg^k n) \quad k > 0$
 $\rightarrow O(n^{\log_b a} \lg^{k+1} n)$

$f(n) = \Omega(n^{\log_b a + \epsilon}) \quad \epsilon > 0$

符合 $af(\frac{n}{b}) \leq cf(n)$ 才能使用
 $\rightarrow O(f(n))$

2.

(a) $T(n) = 4T(n/2) + n$

$\rightarrow n = O(n^{\log_2 4 - \epsilon}) \quad \epsilon = 1$

$T(n) = O(n^2) \#$

(b) $T(n) = 4T(n/2) + n^2$

$\rightarrow n^2 = O(n^{\log_2 4} \lg^k n) \quad k = 0$

$T(n) = O(n^2 \lg n) \#$

(c) $T(n) = 4T(n/2) + n^3$

$\rightarrow n^3 = \Omega(n^{\log_2 4 + \epsilon}) \quad \epsilon = 1$

$4(\frac{n}{2})^3 \leq cn^3, \frac{n^3}{2} \leq cn^3, c = \frac{1}{2}$

$T(n) = O(n^3) \#$

2(b)

$\begin{cases} T(n) = F[1000] \\ T(n) = n \\ F[0] = 0 \\ F[1] = 1 \\ T(n+1) = \end{cases}$

$n-1+1$

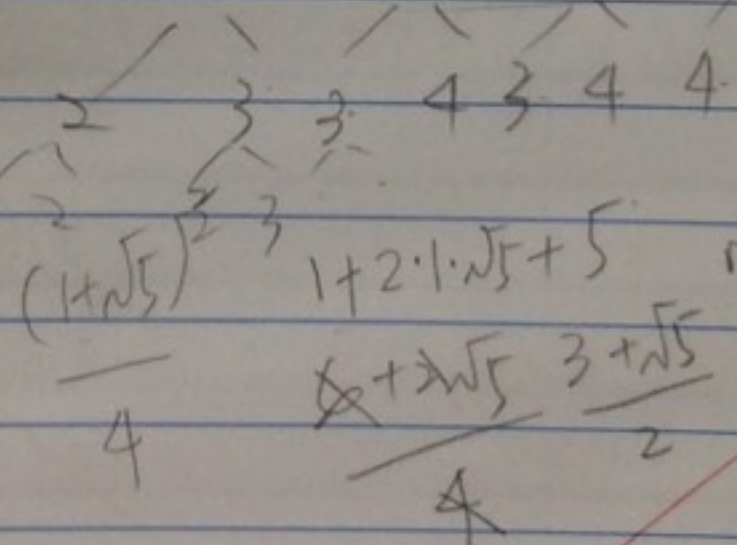
for i from 2 to n
do

$n-1+1$

$F[i] \leftarrow F[i-1] + F[i-2]$

3.

$\begin{cases} 1 & n=1, n=2 \\ T(n-1) + T(n-2) \end{cases}$



$\geq n+5 \rightarrow O(n)$

$t^2 - t - 1 = 0$

$t = \frac{1 \pm \sqrt{5}}{2}$

$A\left(\frac{1+\sqrt{5}}{2}\right)^n + B\left(\frac{1-\sqrt{5}}{2}\right)^n$

(a) $O\left(\frac{1+\sqrt{5}}{2}\right)^n$ worst-case. #

Golden ratio

$$t^2 - t - 1 = 0$$

$$t = \frac{1 \pm \sqrt{5}}{2}$$

$$A\left(\frac{1+\sqrt{5}}{2}\right)^n + B\left(\frac{1-\sqrt{5}}{2}\right)^n \quad (a) \quad \theta\left(\frac{1+\sqrt{5}}{2}\right)^n \text{ worst case. } \#$$

Golden ratio

4.1
(a) $T(n) = \begin{cases} 1, & n=1, n=2 \\ 2T(n-2) + T(n-1) + 1, & n \geq 3 \end{cases}$

x-2
x+1

$$x^2 - x - 2 = 0$$

$$x = 2 \text{ or } -1$$

$$y = 2y + y + 1$$

$$-2y = 1$$

$$y = -\frac{1}{2}$$

$$A(2)^n + B(-1)^n - \frac{1}{2}$$

5,
 $R_L(1, n, k) = \begin{cases} R_L(1, n-1, k) - q_n \cdots q_{n-k+1} \cdot p_{n-k} R_L(1, n-k-1, k), & n > k \\ 1 - q_1 \cdots q_k, & n = k \\ 1, & n < k \end{cases}$

$$R_L(1, n, k) = p_1 R_L(1, n-1, k) + p_2 q(1, n-2, k) + p_3 q^2(1, n-3, k) \\ = p_n q^n R_L(1, n-k-1, k)$$

$$T(1) = 2A - B - \frac{1}{2} = 1 \quad \text{--- (1)}$$

$$+ T(2) = 4A + B - \frac{1}{2} = 1 \quad \text{--- (2)}$$

$$6A - 1 = 2$$

$$6A = 3$$

$$A = \frac{1}{2}$$

1. 代λ①

$$1 - B - \frac{1}{2} = 1$$

$$-B = \frac{1}{2}$$

$$B = -\frac{1}{2}$$

$$T(n) = \frac{2^n}{2} - \frac{(-1)^n}{2} - \frac{1}{2}$$

$$= \frac{2^n - (-1)^n - 1}{2}$$

$$(a) \theta(2^n)$$

$$(b) \theta(n)$$

$$(c) \frac{2^n - (-1)^n - 1}{2} \quad \#$$

b.

$$K=1 \quad 2^0 < n \leq 2^1 \quad X=K+3$$

$$K=2 \quad 2^1 < n \leq 2^2 \quad = \lceil \log_2 \log_2 n \rceil + 3$$

$$K=3 \quad 2^2 < n \leq 2^3$$

$$\rightarrow 2^{K-1} < n \leq 2^K$$

$$K-1 < \log_2 \log_2 n \leq K$$

$$\log_2 \log_2 n \leq K \leq \log_2 \log_2 n + 1$$

(a) $O(\log_2 \log_2 n)$

(b) $\lceil \log_2 \log_2 n \rceil + 3$

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(a) Int Unbiased-Rand-Bit()

```

{
  Int a, b;
  while (true) do
  {
    a ← Biased-Rand-Bit();
    b ← Biased-Rand-Bit();

    if (a ≠ b)
      return 1;
  }
}

```

(b) $\Pr\{\text{1st Interact return 1}\} = P(1-P)$

$\Pr\{\text{1st Interact return 0}\} = (1-P)P$

$\Pr\{\text{1st no return}\} = 1 - 2P(1-P) = \frac{1}{2}$

$\Pr\{\text{1st return 1}\} = \frac{P(1-P)}{1 - \frac{1}{2}} = \frac{P(1-P)}{2P(1-P)} = \frac{1}{2}$

$$P(1-P) + \frac{1}{2}P(1-P) + \frac{1}{2}P(1-P) = \dots$$

$$= \frac{P(1-P)}{1 - \frac{1}{2}} = \frac{P(1-P)}{2P(1-P)} = \frac{1}{2}$$

if (a ≠ b)
return 1;

(c) $i+h \quad 1 \dots n$

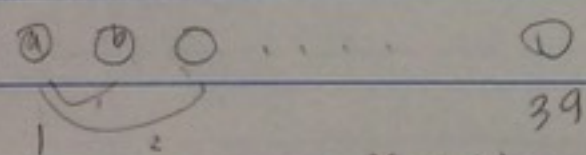
$$E(X) = E(X_1) + E(X_2) + \dots + E(X_n)$$

$$= 1 \times P(P-1) + 1 \times \frac{1}{2}P(P-1) + \frac{1}{4}P(P-1)$$

$$= \frac{1}{1-\frac{1}{2}} = \frac{1}{2P(P-1)} \#$$

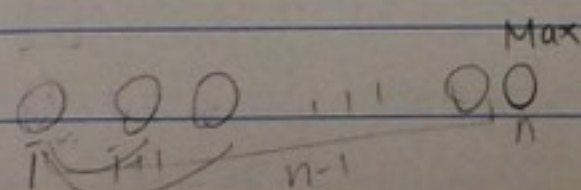
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8.



(a) $38 = k$ if (a > b) swap(a, b)
 $O(n)$

(b)



① 找出 Maximum $n-1$
[1...n-1] 找出最大 $n-2$

$$2n-3 = k$$

$$2 \cdot 39 - 3 = k$$

$$O(n)$$

(c) Maximum $n-1$
 $n-2$
Minimum 1

$$\frac{1 + (n-1) \times n}{2} = \frac{n^2}{2}$$

$$O(n^2)$$