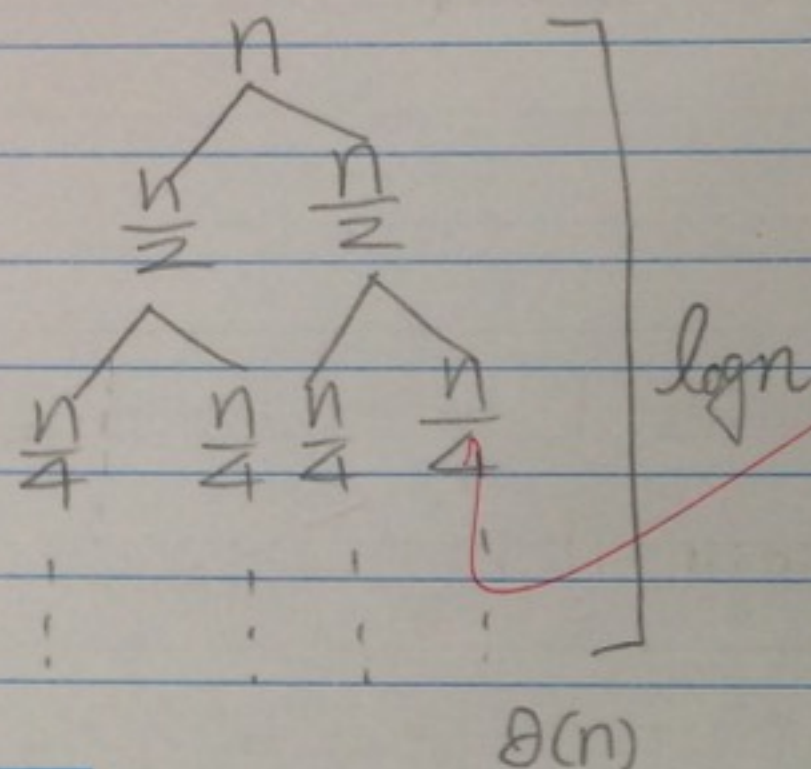


(上列各項，考生須逐項填明，違者該科試卷以零分計)

Please fill out the above information, otherwise the score of this quiz will be zero

①  $A[1, \dots, n]$ 1. If  $n=1$  done2. 遞迴 Sort  $A[1, \dots, \lfloor \frac{n}{2} \rfloor]$  and  $A[\lfloor \frac{n}{2} \rfloor + 1, \dots, n]$ 

3. Merge Sort 2 lists



$$\therefore \Theta(n \log n) \#$$

② (a)  $f(n) = n$ , by CASE 1:  $n^{\log_2 4 - \epsilon} = n^{2 - \epsilon}$ ,  $\epsilon = 1$ 

$$\text{則 } T(n) = \Theta(n^2)$$

(b)  $f(n) = n^2$ , by CASE 2:  $n^{\log_2 4} \lg^k n = n^2 \lg^k n$  令  $k=0$ 

$$\text{則 } T(n) = \Theta(n^2 \lg n) \#$$

(c)  $f(n) = n^3$ , by CASE 3:  $n^{\log_2 4 + \epsilon} = n^{2 + \epsilon}$ , 令  $\epsilon = 1$ , 且  $4f(\frac{n}{2}) \leq c f(n)$ 

$$\text{則 } T(n) = \Theta(n^3) \#$$

$$\begin{aligned}
 \textcircled{3} \quad (1234) \bmod 100 &= (34) \bmod 100 = ((34^2) \bmod 100) \times 34 \bmod 100 \\
 &= ((56) \bmod 100) \times 34 \bmod 100, = 56 \times 34 \bmod 100 \\
 &= 24 \bmod 100.
 \end{aligned}$$



$$\textcircled{4} \quad T(n) = \begin{cases} 1 & n=1 \text{ or } n=2 \\ 2T(n-2) + T(n-1) + 1 & n > 2 \end{cases}$$

$$\text{令 } \lambda^2 = 2 + \lambda \text{ 則 } \lambda = -1 \text{ or } 2$$

$$\beta = 2\beta + \beta + 1, \beta = -\frac{1}{2}$$

$$\text{則 } T(n) = A(2)^n + B(-1)^n - \frac{1}{2}$$

$$\text{令 } T(1) = 2A - B - \frac{1}{2} = 1$$

$$T(2) = 4A + B - \frac{1}{2} = 1$$

$$A = \frac{1}{2}, B = -\frac{1}{2}$$

$$\text{則 } T(n) = \frac{1}{2}(2)^n - \frac{1}{2}(-1)^n - \frac{1}{2}$$

$$(c) = \frac{(2)^n - (-1)^n - 1}{2}$$

$$(b) = \theta(n)$$

$$(a) = \theta(2^n)$$



$$(a) \cdot R_L(1, n, k) = \begin{cases} 1 & n < k \\ 1 - q_1 q_2 \dots q_k & n = k \\ R_L(1, n-1, k) - p_{n-k} q_{n-k+1} \dots q_n R_L(1, n-k-1, k) & n > k \end{cases}$$

$$(b) = R_C(1, n, k) = \begin{cases} 1 & n < k \\ \sum_{i=1}^k q_1 q_2 \dots q_i p_{i+1} R_C(i+1, n+i-1, k) & n \geq k \end{cases}$$

⑥

Randomized Algorithms: 有一 Input 值, 会随机产生不同 output 值

Non-Randomized Algorithms: 有一 Input 值, Output 有可能相同。

⑦

int Biased-RBq ( )

int m, n;

while (true)

{ if (  $\frac{m}{m+n} < \frac{n}{m+n}$  )

return 0;

else return 1;

m ← Unbiased-RBq(m);

n ← Unbiased-RBq(n);

(b)

$$\Pr\{\text{Output 1 1st itera}\} = \frac{n}{m+n} \cdot (1 - \frac{n}{m+n}) = \frac{mn}{m+n}$$

$$\Pr\{\text{Output 0 1st itera}\} = \frac{m}{m+n} \cdot (1 - \frac{m}{m+n}) = \frac{mn}{m+n}$$

$$\Pr\{\text{No output 1st itera}\} = (1 - 2 \frac{mn}{m+n}) = q$$

$$\Pr\{\text{Output 1}\} = \frac{mn}{m+n} + q \left( \frac{mn}{m+n} \right) + q^2 \left( \frac{mn}{m+n} \right) + \dots$$

$$= \frac{\frac{mn}{m+n}}{1 - (1 - 2 \frac{mn}{m+n})} = \frac{1}{2}$$

$$\Pr\{\text{Output 0}\} = \frac{mn}{m+n} + q \left( \frac{mn}{m+n} \right) + q^2 \left( \frac{mn}{m+n} \right) + \dots = \frac{1}{2}$$

(c)

$$E(x) = 1 \times 1 + 1 \times q + 1 \times q^2 + \dots$$

$$= \frac{1}{1-q} = \frac{1}{1 - (1 - 2 \frac{mn}{m+n})} = \frac{1}{2 \frac{mn}{m+n}} = \theta\left(\frac{1}{2 \frac{mn}{m+n}}\right)$$

⑧ 令 k = 执行次数

$$(a) \quad k=1 \quad 2 < n \leq 2$$

$$(b) \quad k=3+x$$



(c)

$$E(x) = |x| + |xq| + |xq^2| + \dots$$

$$= \frac{1}{1-q} = \frac{1}{1-(1-2\frac{mn}{m+n})} = \frac{1}{2\frac{mn}{m+n}} = \theta\left(\frac{1}{2\frac{mn}{m+n}}\right)$$

⑧ 令  $k$  = 執行次數

a)

$$k=1 \quad 2^0 < n \leq 2^1$$

$$(b) \quad \text{又 } k = 3 + X$$

$$k=2 \quad 2^1 < n \leq 2^2$$

$$n = 3 + \lceil \log_2 \log_2 n \rceil$$

$$k=n \quad 2^{k-1} < n \leq 2^k$$

$$\downarrow \times \log_2 \log_2$$

$$k-1 < \log_2 \log_2 n \leq k$$

$$\therefore \theta(\log_2 \log_2 n)$$

⑨ SELECT( $\bar{i}, n$ )

$\theta(n)$ : 1. Divide  $n$  elements into group of 5 找出中間數  $x$ ,

$\theta(\frac{n}{5})$ : 2. 找出所有中間數的中間數  $\lceil \frac{n}{5} \rceil$  (用遞迴)

$\theta(n)$ : 3. partition around the  $x$ , let's  $k = \text{rank}(x)$

$\theta(\frac{3}{4}n)$ : 4. if  $\bar{i} = k$  return  $x$

else if  $\bar{i} < k$  then select  $\bar{i}$ -th smallest in lower part,

else select  $(\bar{i} - k)$ -th smallest in upper part.

$$\therefore O(n) = T(\frac{n}{5}) + T(\frac{3}{4}n) + \theta(n) \stackrel{!}{=} \theta(n)$$