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CS143  
HW 5

1)

To read all Tuples of R,

$$750,000/500 = 1,500 \text{ blocks}$$

Since for every tuple in R there are roughly 5 tuples in S, to read the tuples in S

$$750,000 * 5 = 3,750,000 \text{ blocks}$$

Total Block read =

$$1,500 \text{ blocks of R} + 3,750,000 \text{ blocks of S} = 3,751,500 \text{ blocks}$$

The new plan requires the same amount of R-block read

However, it requires an extra S-block read for

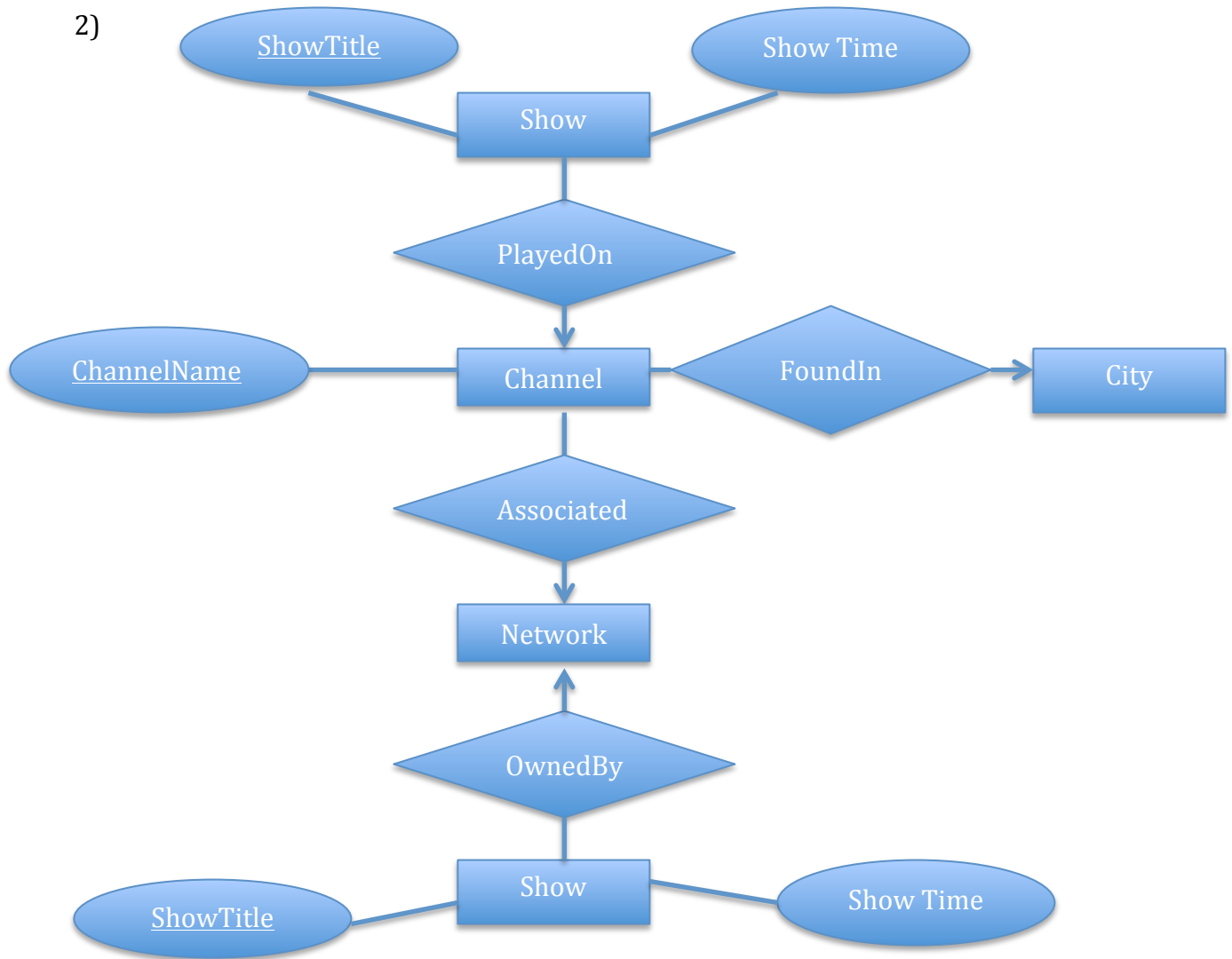
$$750,000 * 5000 / 100 \text{ sequential} = 37,500,000 \text{ blocks}$$

Total Block read =

$$37,500,000 + 1,500 = 37,501,500 \text{ blocks}$$

The second plan requires 10 times more block access than the first plan. However, if the cost of random access is much more ( $>10$ ) times than sequential access, the second plan will require less time.

2)



PlayedOn(ShowTitle, Channel)  
Associated(ChannelName, Network)  
FoundIn(ChannelName, City)  
OwnedBy(ShowTitle, Network)

3)

Parts(number)  
Assembly(number, cost)  
ComposedOf(number, cost, quality)

4)

$R1(A,B,C,F) \cap R2(A,D,E) = A$

$A \rightarrow BC \rightarrow \underline{CD} \rightarrow \underline{E} \rightarrow \underline{A}$

$A \rightarrow ADE$

$R1(A,B,C,F) \cap R2(A,D,E) \rightarrow R2(A,D,E)$

It is a lossless decomposition

5)

$A \rightarrow B$

$C \rightarrow A$

This is because B is always constant and C is the most varied.

6)

a)  $E \rightarrow A, A \rightarrow BC, BC \rightarrow CD, CD \rightarrow E$

$\{E\}^+ = \{ABCDE\}$

E is a key

b)  $BC \rightarrow CD, CD \rightarrow E, E \rightarrow A, A \rightarrow BC$

$\{BC\}^+ = \{ABCDE\}$

BC is a key

7)

NOT BCNF

$\{A\}^+ = \{BCDE\}$  Not BCNF

$\{B\}^+ = \{D\}$  Not BCNF

$\{C\}^+ = \{E\}$  Not BCNF

Normalized

$R1(A,B,C)$

$R2(B,D)$

$R3(C,E)$