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/* CS261- Assignment 3 - Part 2: Comparison Questions*/
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- 1) Which of the implementations uses more memory? Explain why.
- * The linked list implementation used more memory than the dynamic array implementation. The linked list began getting memory feedback at n = 16,000 whereas, the dynamic array didn't get its first memory feedback until n = 64,000.
- * The linked list utilized more memory than the dynamic array because for one, the linked list had a greater intial overhead due to the pointers that are member variables of the struct LinkedList and struct DLink's. So if as n increases, more elements are created that require extra memory for their own respective member pointers. This would also explain why both implementations got memory feedback at differnt n's. The linked list would at n = 16,000 would have to have memory for the several member pointers that each of the 16,000 link's would require. The dynamic array would not require this extra allottment of memory and so at n = 16,000 there would only need to be memory for the elements of the array and their 1 TYPE member variable. Even if the linked list only had 1 pointer member, for n elements, you need extra memory for n pointers.
- 2) Which of the implementations is the fastest? Explain why.
 - * The dynamic array implementatin had the fastest run times.
- * The dynamic array's runtime advantage can be attributed to several variables. First, the linked list has guaranteed O(1) on insertion and removal (depending on whether the list is single or doubly linked), the dynamic array has an amortized O(1)+ on insertion on insertion and removal operations by virtue of the nature of the dynamic array as a structure. However, the dynamic array's allocations and deallocations of memory do not occur every insertion or removal operation, only when size = capacity or when size = 1/2 capacity (or some variation). Second, due to the nature of the linked list as a structure, the elements of the linked list do not exist contiguously in memory as they do in the dynamic array. So when performing an insertion or removal operation on a linked list, the positioning of the next element is not a constant distance, and so execution times can be slowed. The dynamic array has the benefit of a cache unit arrangement which optimizes the execution of insertion and removal operations, and the dynamic array's allocations and deallocations of memory occur less frequently than the linked list, resulting in an amortized run time of O(1)+. So the relational constraints that make up the dynamic array abstraction, result in faster run times in relation to linked list structures.
- 3) Would you expect anything to change if the loop performed remove() instead of contains()? If so, what?
- * Using removeList() in place of containsList() I would expect an increase in the run time of the linked list because in removeList(), if the value of the element in question (index i in a loop) is equal to the value being removed, then the next and prev pointers of the list need to be reassigned to the their logically respective neighbors (shifted) to represent the removal of an element. Memory would also need to be allocated for a struct DLink* temp that is used to reattach the links post removal and then this temp must be freed to remove the value to be removed. The linked list has the benefit not requiring a nested loop within its containsList(), however, memory must be allocated to a temp pointer (which

also must be freed) and member pointers of neighbor links must be reassigned to represent the reattachment of the linked list post link removal. For these reasons I would expect the swapping of the containsList() and the removeList() to increase the run time of this linked list.

* Using removeList() in place of containsList() I would also expect an increase in the run time of the dynamic array because in removeList(), if the value of the element in question (the element used as an argument of the conditional) another loop runs within the initial n outside loop. This inner loop runs from the index of the element with the matching value of the value to be removed, until the index of the size of the dynamic array. This inner loop shifts all values from the removed element til the last element in the dynamic array. In the containsList() the was a loop that ran n times, but even if the conditional within that loop was satisfied, there were no more loops and there was no shifting of elements and values. So by swapping removeList() for containsList() the dynamic array's runtime increases due to a nested loop within the function and the required shifting of values when a value to be removed is found within the dynamic array.





