



Separation of Data flow and Control flow in Reconfigurable Multi-core SoCs using the Gannet Service-based Architecture

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- Overview
 - Task control in NoC-based SoCs
 - Gannet system
 - Gannet language
 - Separation of data flow and control flow



Overview

- We propose a novel approach to **separating control flow from data flow** in NoC-based SoCs consisting of multiple heterogeneous reconfigurable IP cores.
- This mechanism
 - enables full data path control by an embedded microcontroller
 - whilst avoiding the potential memory bandwidth bottleneck and
 - without requiring centralised control over the NoC.



Overview

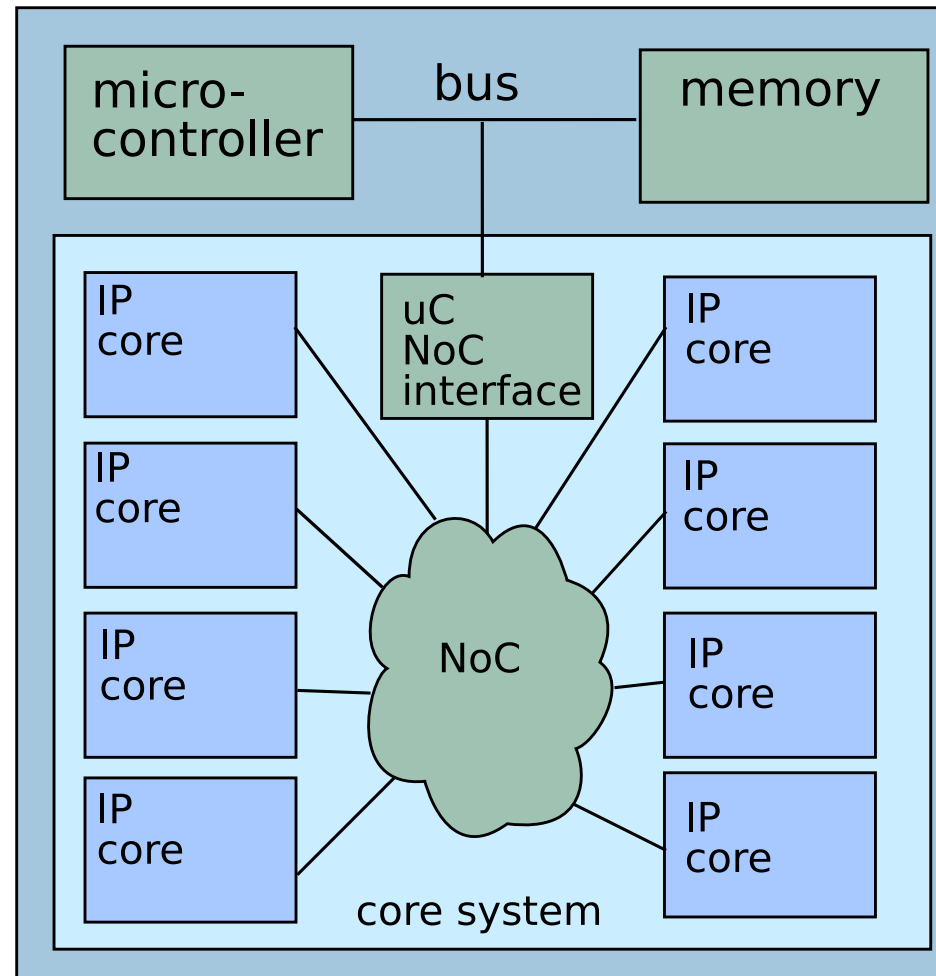
- We assume a generic SoC where
 - data is processed by IP cores interacting through a NoC
 - control structures are implemented on a microcontroller.
- The proposed mechanism employs a service-based SoC architecture (the **Gannet** architecture) where
 - the control services are implemented using a Virtual Machine
 - IP cores acquire service behaviour through the use of a generic data marshalling and interfacing circuit.



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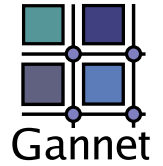
Task-level reconfigurability

- NoC-based SoC with embedded microprocessor:



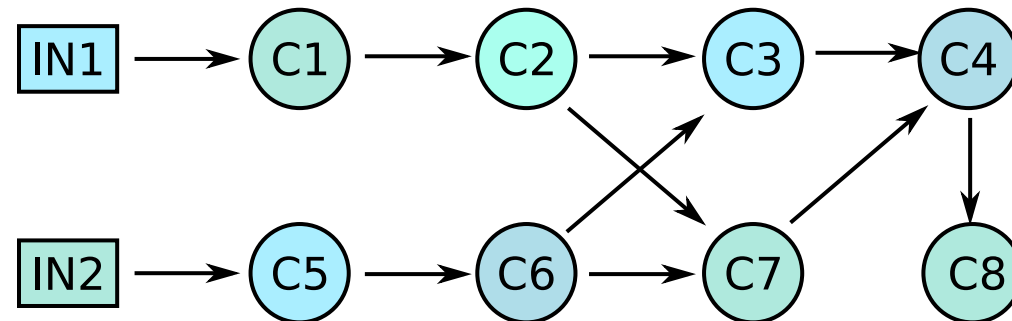


Task control in NoC-based SoCs

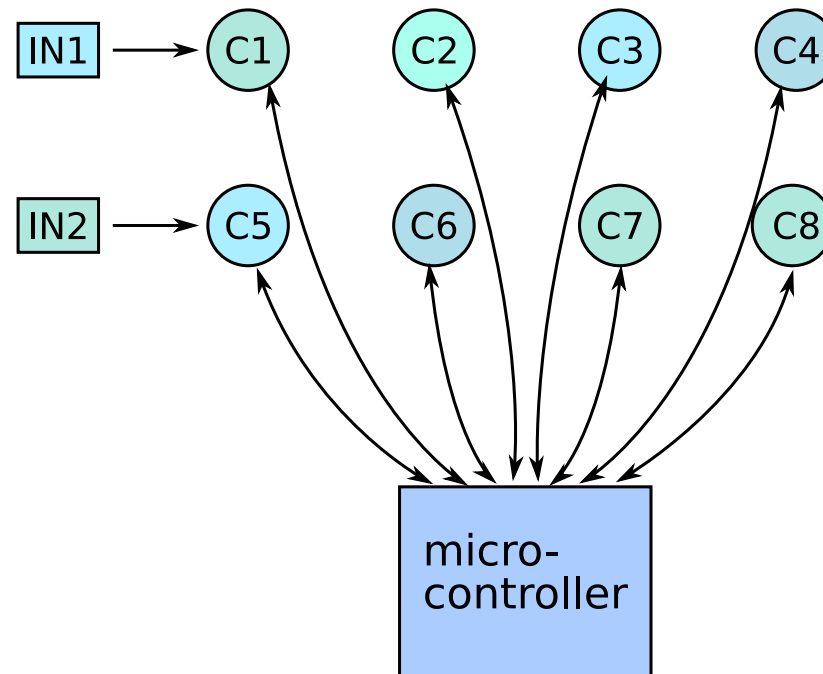


- SoCs in general use an embedded microprocessor for control.
- Conventional way of controlling hardware blocks using an embedded microprocessor: memory-mapped IO+ interrupts.
- In a NoC-based SoC, the microprocessor interacts with a NoC transceiver and transfers data as NoC packets \Rightarrow
 - efficient data transmission;
 - considerably reduction of required number of interrupts;
 - no significant operational difference with bus-based mechanism.

- Non-task-level reconfigurable system:
 - microcontroller only sends control or configuration information to each core;
 - all data can flow between the cores.



- Task-level reconfigurable system:
 - data paths are determined at run time by a program running on microcontroller;
 - all data pass via the microcontroller.



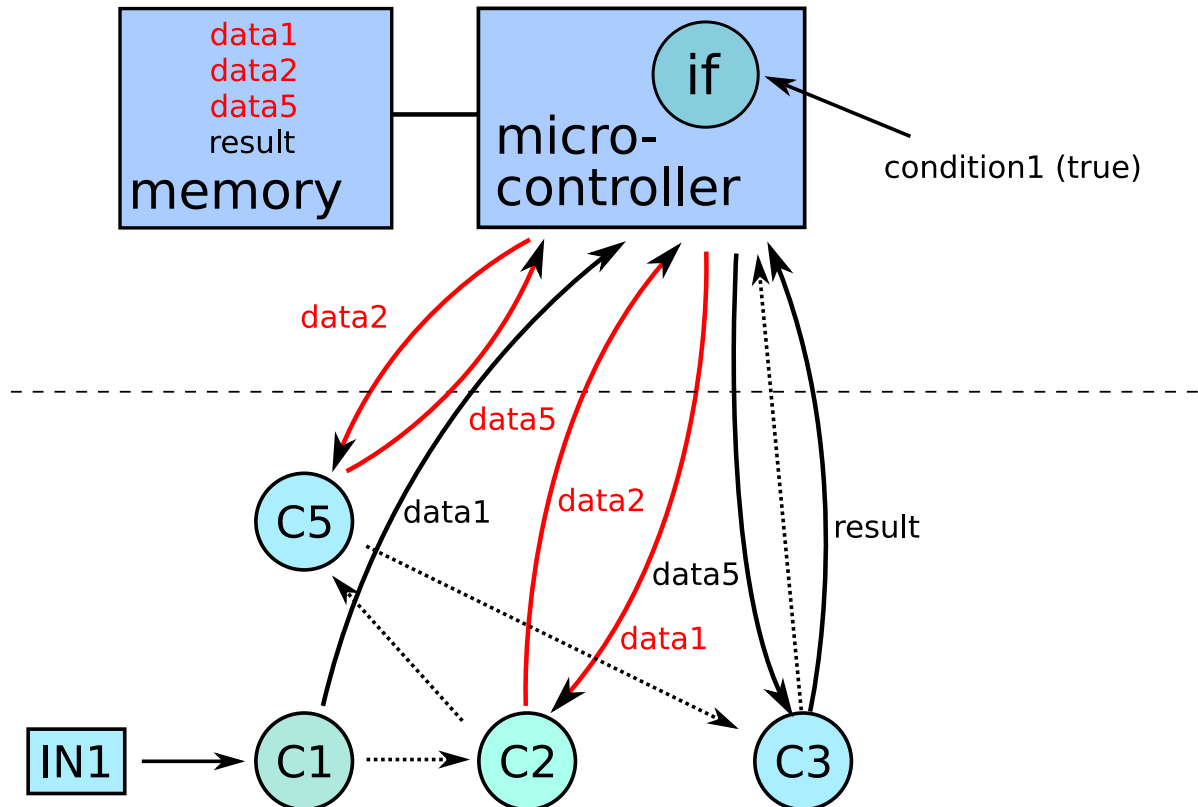


Example

```
Data* NoC_TRX(CoreAddres&,...); // variadic function prototype
/* variable declarations omitted */
data1=NoC_TRX(C1,IN1);
if (condition1) {
    Data* data2=NoC_TRX(C2,data1);
    data5=NoC_TRX(C5,data2);
} else {
    Data* data2=NoC_TRX(C5,data1);
    data5=NoC_TRX(C2,data2);
}
result=NoC_TRX(C3,data5);
```



Example





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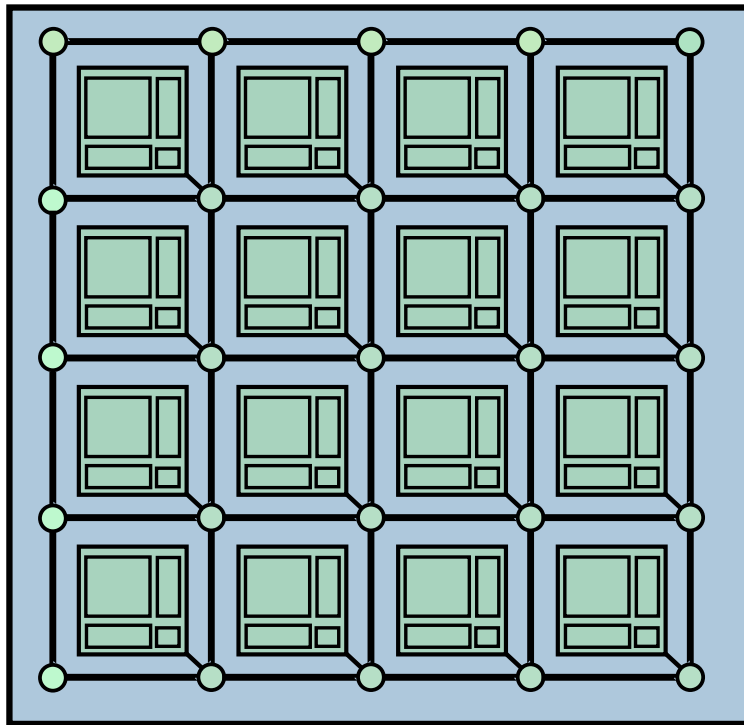


Gannet architecture

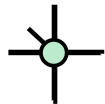
- A **service-based** architecture for **very large SoCs**:
 - a collection of processing cores (HW/SW).
 - each core offers a a specific **service**.
 - tasks are defined by the interaction pattern of the services.
- Task-level reconfigurability
 - task description programs, configurable at run time
- High abstraction-level design
 - single program governs behaviour of complete system



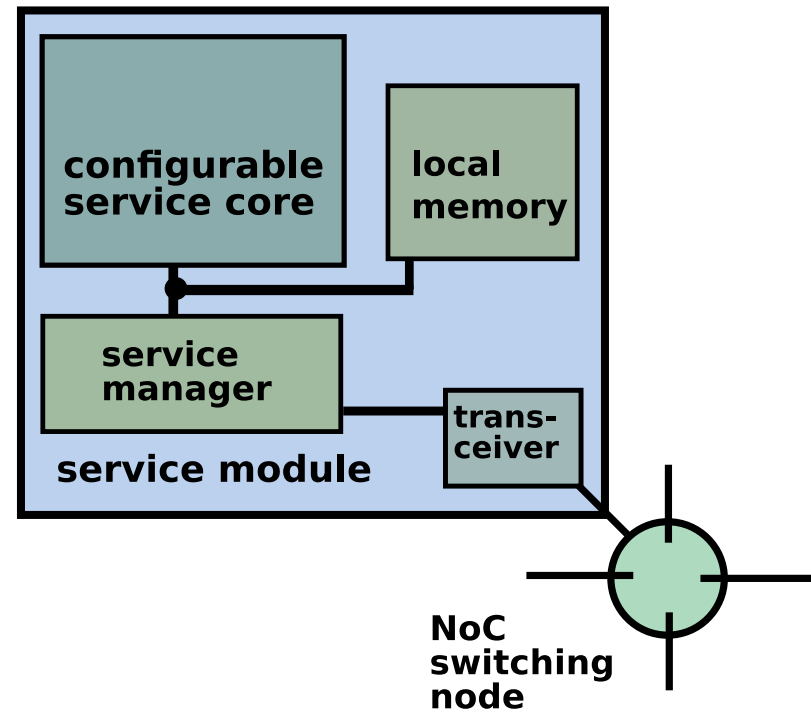
Gannet architecture



tile with
IP core



NoC switch





Gannet system operation

- The Gannet machine is a distributed computing system where every computational node **consumes packets** and **produces packets** and can store state information between transactions.
- We denote a Gannet packet as $p(\textit{Type}, \textit{To}, \textit{Ret}, \textit{Id}; \textit{Payload})$
- The semantics of a Gannet **service** (computational node) can be described in terms of
 - the **task code**
 - the internal state
 - the **result packet(s)** produced by the task

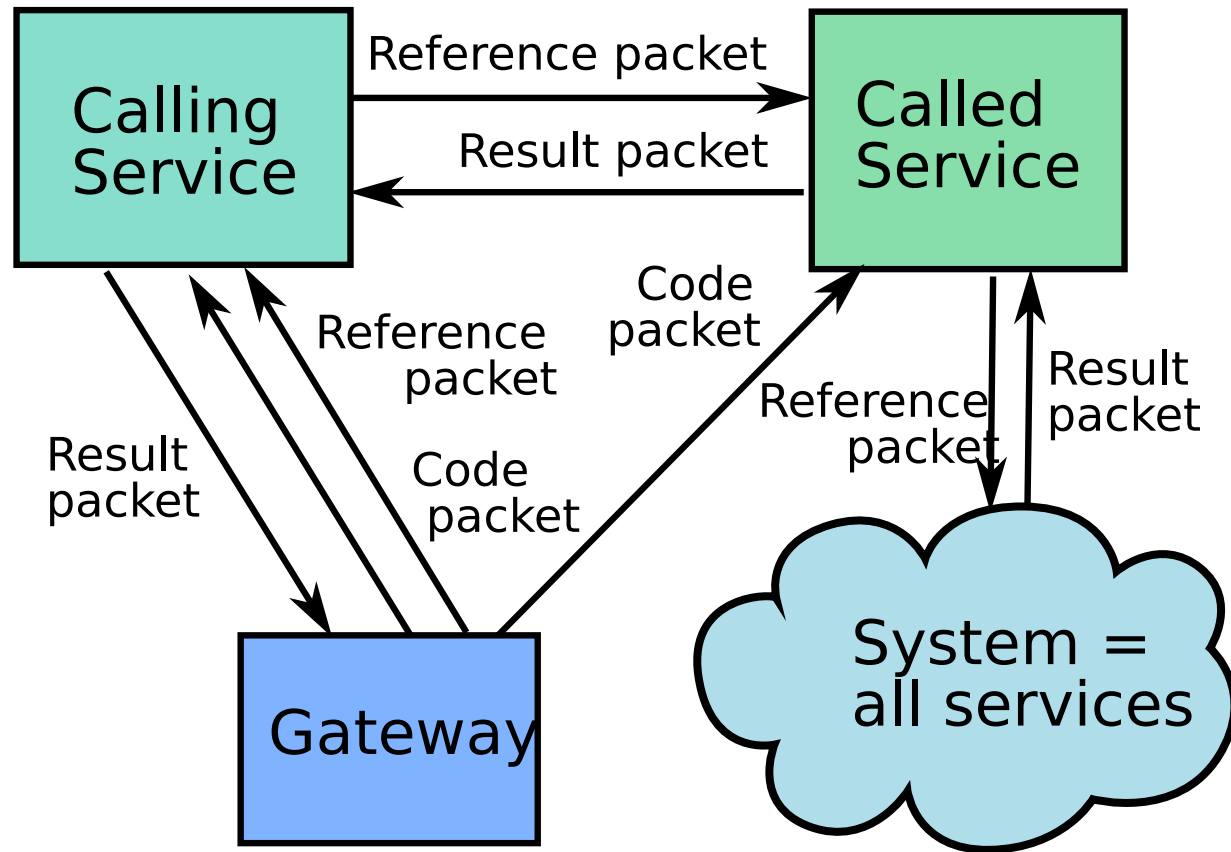


Gannet system operation

1. Service S_i receives a **code** packet $p(\text{Code}, S_i, S_j, R_{task}; task)$ where $task = (S_i a_1 \dots a_n)$. The task is stored with reference R_{task} . Service S_i is in $state_i$.
2. Service S_i receives a task **reference** packet $p(Ref, S_i, S_j, R_{id}; R_{task})$
3. The service activates the task referenced by R_{task} : $(S_i a_1 \dots a_n)$.
This results in evaluation of the arguments $a_1 \dots a_n$:
4. The service produces a result packet $p(\text{Type}_i, S_j, S_i, R_{id}; Result_i)$ and the state changes to $state_i'$.
5. This packet is sent to S_j where $Result_i$ is stored in a location referenced by R_{id} .



Gannet system operation





Control services in Gannet

- Any run-time reconfigurable system requires **control constructs** to be effective.
- In Gannet, these constructs (if/then, functions, blocks, variables, ...) are provided by **services**.
- Such **control services** can be efficiently implemented on an embedded microcontroller.
- Interleaving the services provided by the HW cores with control services can cause bottleneck due to memory bandwidth.



Control services in Gannet

- Ideally, the microcontroller would only exchange control information with the cores.
- technically not impossible to realise using compiled code but would require
 - a language with functional characteristics (no side-effects, undetermined execution order, laziness, concurrency)
 - access to absolute memory addresses of the data structures
- program would need to contain a JIT compiler to create bytecode for the service managers at run time.

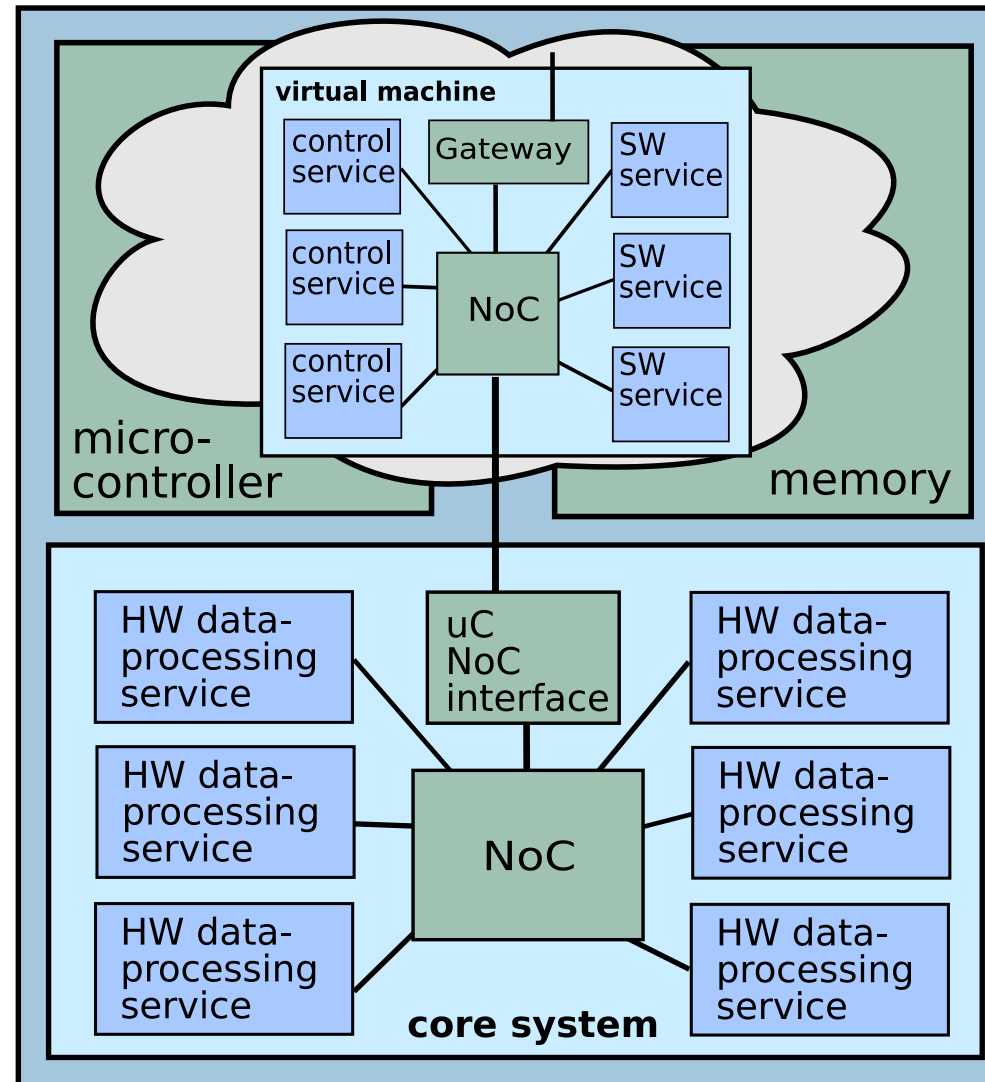


The Gannet Virtual Machine

- A Virtual Machine (VM) which interacts with the hardware service managers:
 - software implementation of the service managers, control service cores and a 'virtual NoC'
 - small, portable C++ application
 - runs byte-compiled programs in the Gannet language
 - same bytecode used by VM and HW



The Gannet system





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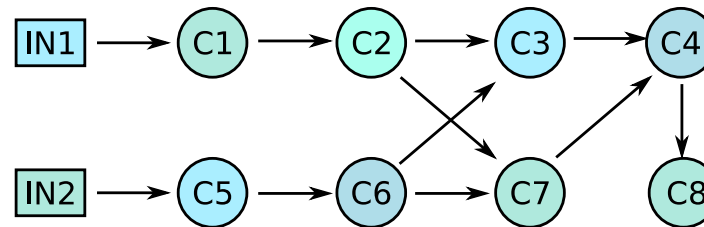
Gannet language

- The “assembly” language to program the Gannet system
- Intended as compilation target, not HLL
- A functional language, every service is mapped to an opaque function
- Gannet is a distributed machine for running this language
- Service = service manager + service core
- Service cores can be implemented in HW or SW



Gannet language syntax

■ Previous example in S-expressions syntax:



```
(S8 (S4  
  (S3  
    (S2 (S1 IN1))  
    (S6 (S5 IN2))  
  )  
  (S7  
    (S6 (S5 IN2))  
    (S2 (S1 IN1))  
  )  
) ) )
```




Gannet language syntax

■ Example with control services (factorial):

```
(let
  (assign 'fact
    (lambda 'n 'a 'f
      '(if (< n '2)
        'n
        '(apply f (- n '1) (* a n) 'f)
      )))
  (apply fact '4 '1 'fact)
)
```



Gannet language properties

- Some key properties of the Gannet language:
 - the evaluation order is unspecified
 - eager by default but lazy evaluation is possible
 - no side effects across services
 - updates of variables are atomic (no race conditions)
- These properties
 - make the language fully concurrent (maximise parallelism)
 - and enable separation of control flow from data flow



■ Unspecified execution order:

- In a given function call it is not possible to predict the evaluation order of the arguments.
- In practice, all arguments are evaluated in parallel; call blocks until all arguments are ready.

```
(let  
  (assign 'a (S1 ...))  
  (assign 'b (S2 ...))  
  (S3 ... b ...  
    (S4 ... a ...)... )  
)
```



Gannet language properties

■ Lazy evaluation:

- By default, Gannet is **eager**, i.e. it always evaluates all arguments before passing them on to the service core.
- It should be possible to evaluate arguments at need (“**lazy**”).
- Laziness is expressed by prefixing an expression or symbol with a single quote:

```
(assign 'a (S1 ...))
```

- Quoting causes the evaluation of the symbol `a` to **deferred** to the service core.



Language properties

- **No side effects across services:**

- A call to a given service should not result in a modification of the state of the rest of the system.

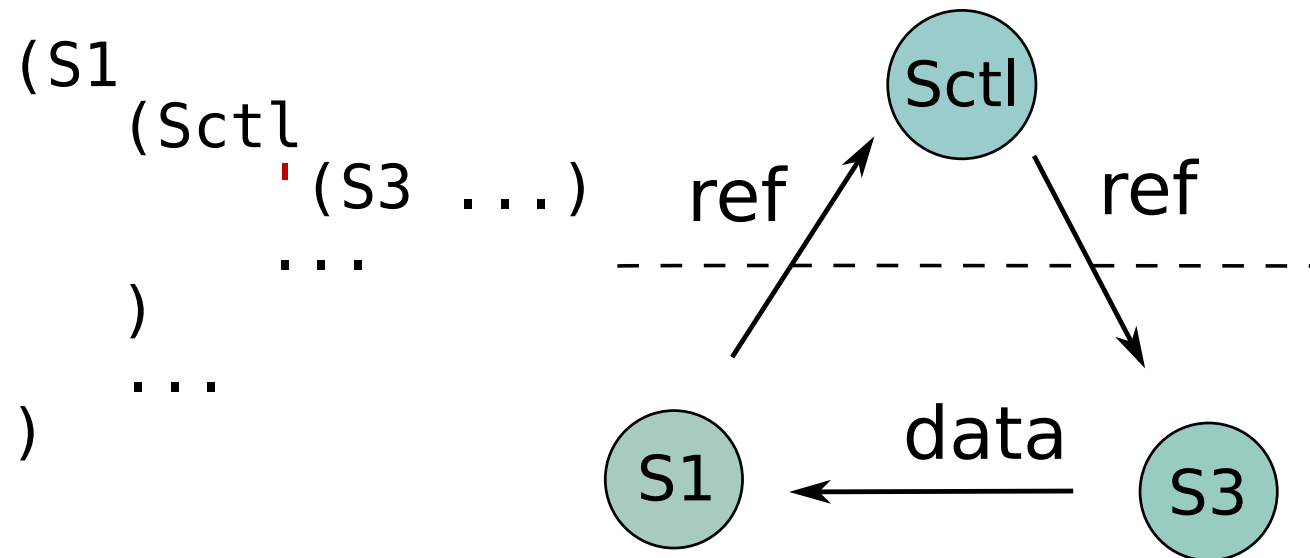
- **Updates of variables are atomic:**

- No race conditions if several services simultaneously try to modify shared data.
- The service manager processes all task requests in FIFO order.
- Not possible to update an unassigned variable.



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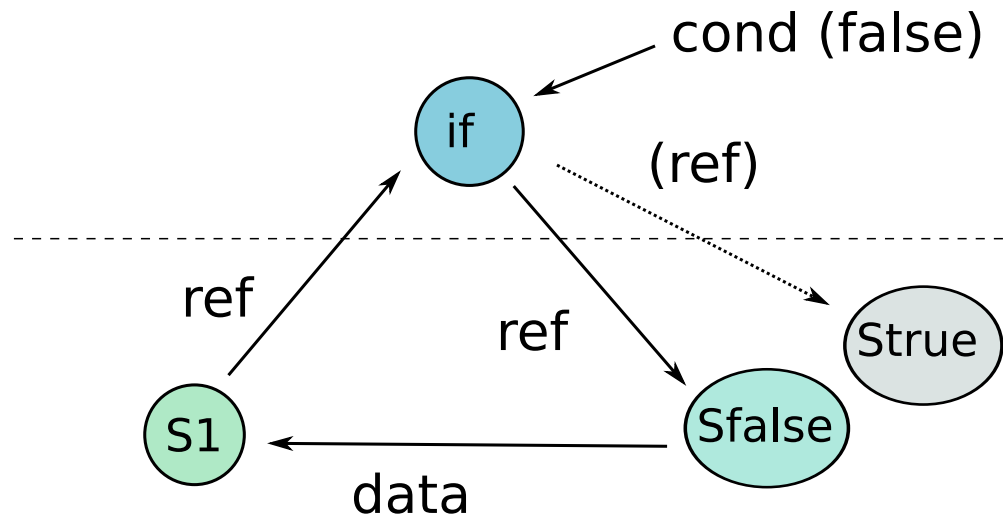
The mechanism for separation of control and data flows: **deferred evaluation** and **redirection**:



S_1 sends $p(Ref, S_{ctl}, S_1, R'_1; R_{ctl})$; S_{ctl} sends $p(Ref, S_3, S_1, R'_1; R_3)$

Conditional branching

```
(S1
  (if cond
    '(Strue ...)
    '(Sfalse ...))
  ...
)
```



S_{if} sends $p(Ref, S_{false}, S_1, R'_1; R_{false})$

Conditional branching

■ Revisiting the earlier example:

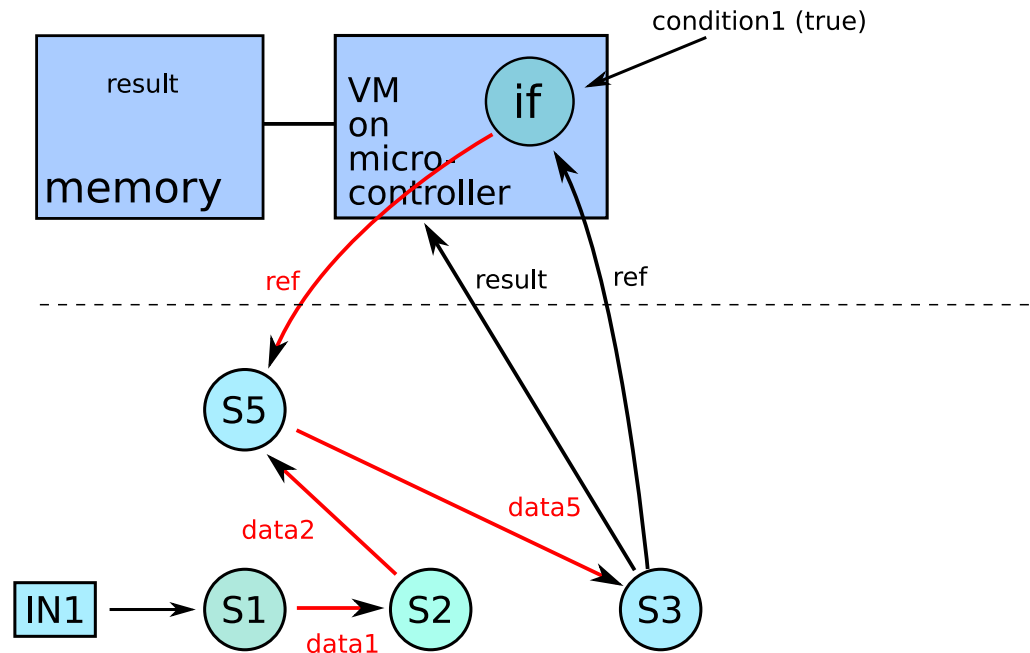
(S3

(if condition1

'(S5 (S2 (S1 IN1)))

'(S2 (S5 (s1 IN1)))

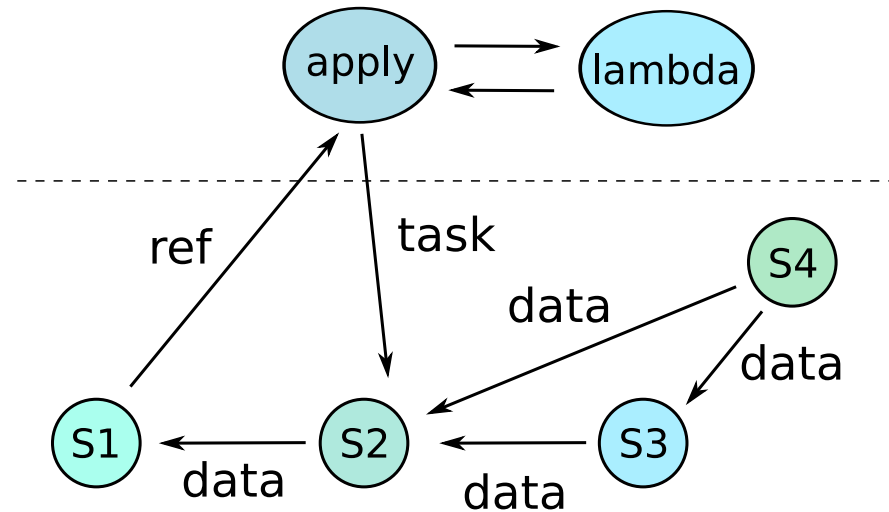
)
)



S_{if} sends $p(Ref, S_5, S_3, R'_3; R_5)$

Function definition and application

```
(S1
  (apply
    (lambda 'x
      '(S2 (S3 ... x ...) ... x ...)
    )
    '(S4 ...)
  )
  ...
)
```

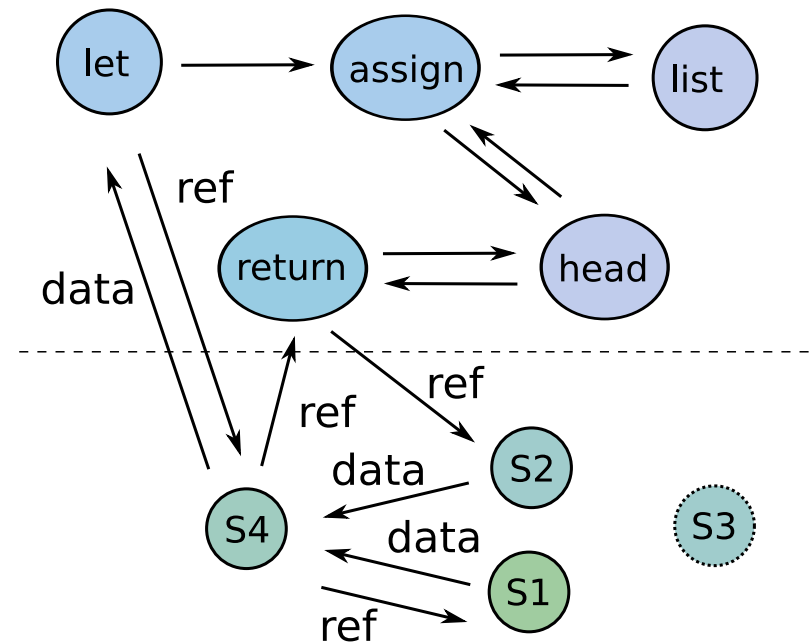


$$S_1((\lambda x \rightarrow S_2(S_3(\dots, x, \dots), \dots, x, \dots))S_4(\dots), \dots)$$

Lists

- `list`: list constructor
- `head`: first element of the list
- `return`: unquotes its argument

```
(let  
  (assign 'l (list '(S2 ...) '(S3 ...)))  
  (S4 (return (head l)) (S1 ...))  
)
```



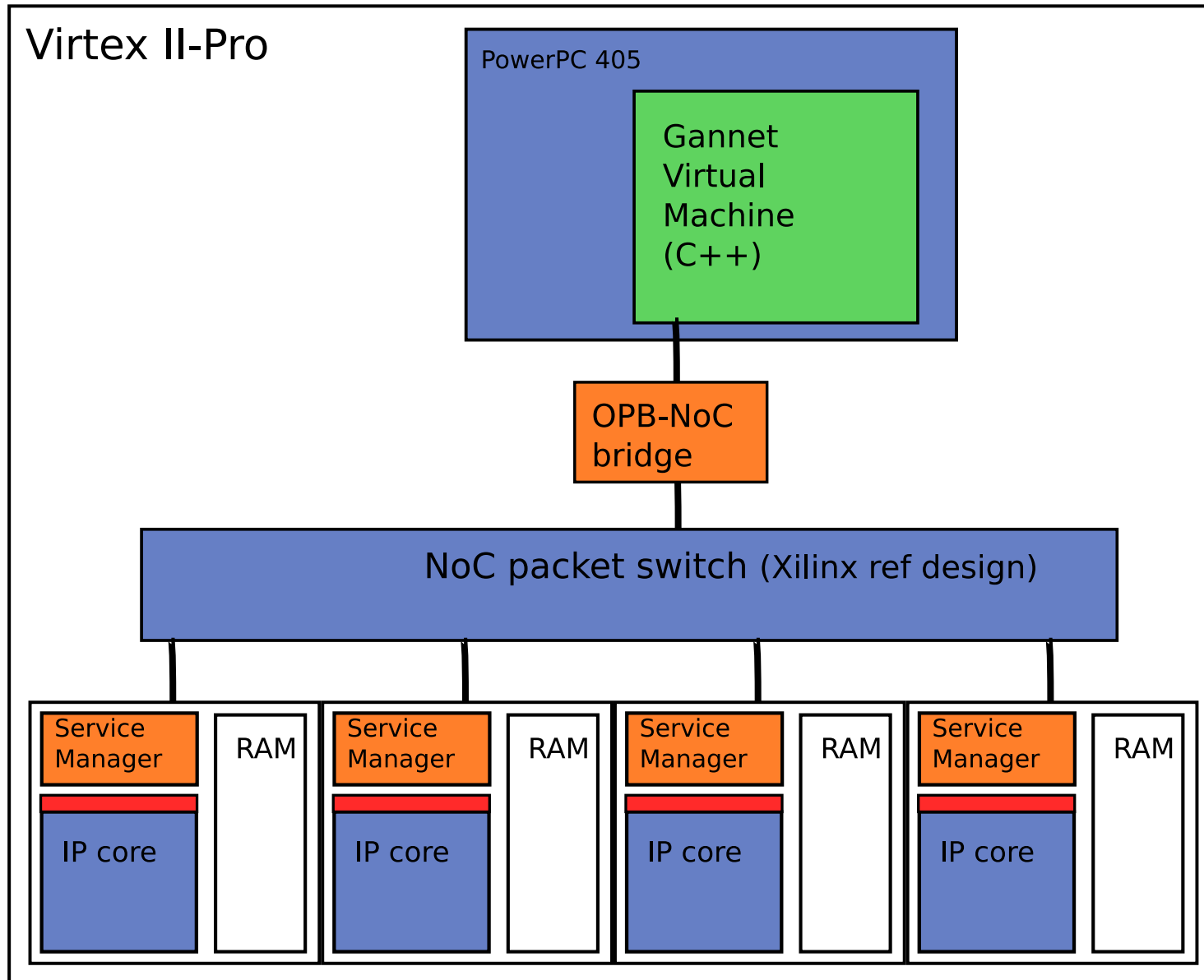


Conclusion

- Gannet: a service-based SoC architecture for high-level design of reconfigurable heterogeneous multi-core SoCs.
- Alleviate bottleneck resulting from memory bandwidth limitation:
mechanism for the **separation of control flow and data flow**
based on **deferred evaluation** and **packet redirection**.
- Gannet system
 - provides full control over data paths in multi-core SoC;
 - provides full concurrency;
 - ensures that data can flow directly between the cores.



Status





Status

