

Module 02:

"Blocking Synchronization"



TEKNOLOGISK
INSTITUT

Agenda

- ▶ **Introducing Synchronization**
- ▶ Synchronization by Locking
- ▶ Best Practices for Locking
- ▶ More Blocking Thread Synchronization
- ▶ Cross-Process Synchronization



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The Need for Synchronization

- ▶ Processor and operating system schedule threads in and out repeatedly
 - Thread context switch can occur at any time
 - Even in the middle of assignments and increments etc.
- ▶ Hence computations need to be computationally safe
 - Some operations must be performed indivisibly!
 - Race conditions should be avoided
- ▶ Basically three solutions
 - **Blocking synchronization** of access to critical regions of code
 - Signaling between threads
 - Nonblocking synchronization

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Mutual Exclusion using Monitors

- ▶ The **Monitor** class is a light-weight mutual exclusion mechanism for use within a single process
 - **Monitor.Enter()**
 - **Monitor.Exit()**

```
object syncObject = new object();  
...  
  
Monitor.Enter( syncObject);  
_counter++;  
Monitor.Exit( syncObject);
```

- ▶ What if we forget to exit?
- ▶ What about exceptions...?

The C# **lock** Keyword

- ▶ The **lock** keyword in C# is based on **Monitor** and **try-finally**

```
object syncObject = new object();  
...  
  
lock( syncObject )  
{  
    _counter++;  
}
```

- ▶ Note: **lock** can only lock reference types...! Why?

Which Synchronization Object?

- ▶ Always choose a reference type instance
- ▶ Best practice is to choose independent, private object
- ▶ Might even give it a descriptive name

```
object _counterAccessSyncObject = new object();  
...  
  
lock( counterAccessSyncObject )  
{  
    _counter++;  
}
```

- ▶ Is the **this** reference a good choice? **Not really!**

Access to Static Members

- ▶ For exclusive access to static members of some type, convention is to use its Type object

```
lock( typeof(Resource) )  
{  
    Counter++;  
}
```

- ▶ Alternatively, create a static synchronization object
 - **typeof(Resource)** suffers similar caveats as **this**

Variations: Monitor.TryEnter()

- ▶ Enter with a timeout
 - Monitor.TryEnter()

```
bool wasAcquired = Monitor.TryEnter( _counterAccessSyncObject );  
if( wasAcquired )  
{  
    _counter++;  
    Monitor.Exit( _counterAccessSyncObject );  
}
```

Variations: Lock Taken Overloads

- ▶ Extreme subtleties lead to more overloads in C# 4.0

```
bool wasLockTaken = false;
try
{
    Monitor.Enter(_counterAccessSyncObject, ref wasLockTaken);
    _counter++;
}
finally
{
    if (wasLockTaken)
    {
        Monitor.Exit(_counterAccessSyncObject);
    }
}
```

- ▶ This is how lock is implemented internally in C# 4

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Locking is... well... Subtle!



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Locking is just Convention!

- ▶ Remember that locks are based on convention
 - They work only in everybody plays along nicely...!

- ▶ Locking
 - Requires discipline from everybody
 - Is easy to forget
 - Is hard to detect that it is forgotten
 - Is just hard to get 100% right!

- Is... Subtle! 😊

Locking Allows Reentrancy

- ▶ Same thread can acquire the same lock **multiple** times
 - Only blocks on the initial attempt to acquire lock

```
lock (_counterAccessSyncObject)
{
    ...
    lock (_counterAccessSyncObject)
    {
        _counter++;
    }
}
```

- ▶ Try to avoid doing this too extensively

Accessing Multiple Resources

- ▶ Use one or more locks for multiple resources

```
lock (_counterFromAccessSyncObject)
{
    lock (_counterToAccessSyncObject)
    {
        _counterFrom--;
        _counterTo++;
    }
}
```

- ▶ **Beware of granularity!**
 - Fewer locks a.k.a. "*Coarse-grained*" => Performance Hit?
 - More locks a.k.a. "*Fine-grained*" => Deadlock Risk?

Deadlocks

- ▶ When
 - Thread 1 has acquired **Resource A** and waits for **Resource B**
 - Thread 2 has acquired **Resource B** and waits for **Resource A**then a deadlock has occurred..!

- ▶ Deadlocks
 - might not occur deterministically
 - cannot be detected automatically by humans or compiler
 - are hard to find and debug!

- ▶ "Livelocks" also exist and are equally painful (but rare)

Lock-Levelling

- ▶ Use strict locking discipline called "Lock-Levelling"
 - Assign some fictitious number to each resource
 - Ensure that any thread always **only locks** a lock with a **higher** number than **any** lock it already holds



- ▶ Alternative is nonblocking synchronization later

Best Practices for Locking

- ▶ Lock access to *any* (writeable,) shared fields!
- ▶ Carefully consider your granularity of locking
- ▶ Use lock-levelling and document the levels in code
- ▶ Never lock when invoking blocking methods
 - E.g. calling external WCF services

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Reader/Writer Locks

- ▶ If there are many readers and only occasional writers, the **ReaderWriterLockSlim** might be more performant

```
try
{
    _rwLock.EnterReadLock();
    Console.WriteLine(_number);
}
finally
{
    _rwLock.ExitReadLock();
}
```

```
try
{
    _rwLock.EnterWriteLock();
    _number = 87;
}
finally
{
    _rwLock.ExitWriteLock();
}
```

- ▶ A write lock is universally exclusive.
- ▶ A read lock is compatible with other read locks.
- ▶ For advanced scenarios read locks may be upgraded to a write lock and vice versa downgraded



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Semaphores

- ▶ Semaphores limit the count of concurrency to a resource
 - Semaphore(1) ~ Monitor
 - Note: Semaphores have no record of owner threads!

```
_semaphore = new SemaphoreSlim(3);  
...  
_semaphore.Wait();  
...  
_semaphore.Release();
```

- ▶ **SemaphoreSlim**
 - Lightweight .NET 4.0 version of **Semaphore**
 - Has asynchronous features
 - Local-only!

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Mutexes

- ▶ A mutex is a cross-process version of Monitor
 - Uses Windows Kernel object in OS

```
_mutex = new Mutex(false, "MyResourceMutex");  
...  
_mutex.WaitOne();  
...  
_mutex.ReleaseMutex();
```

- ▶ Can be both local and cross-process, e.g.
 - Ensure mutually exclusive access to machine-wide resource
 - Ensure at most one instance of application is running

Semaphore

▶ Semaphore

- Uses Windows Kernel object in OS
- Can be both local **and cross-process**

▶ SemaphoreSlim

- Faster and better than **Semaphore** when local
- **Local-only!**

Summary

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