COMP603: MIDTERM I

NAME: _____

		utes. Read each question carefully. Write legibly and check your work. aptops are allowed. Good luck!
		1. Short Definitions
Corr	ectly define 8 of the	following terms for full credit. Correctly define all for extra credit.
(1)	String	sequence of characters/symbols.
(2)	Language	set of strings.
(3)	Compiler	translates a source language to a target language.
(4)	Interpreter	executes source code.
(5)	Bootstrapping	the process of writing a compiler in its own language
	by first writing	a compiler for the language in an existing language.
(6)	Visitor	design pattern for tree traversal.
(7)	Nondeterminism	having a choice about which state to transition to.
(8)	Ambiguity	more than one parse of a string is possible.
(9)	First set	the set of terminals appearing first in any string derived
	from a nontermina	1

(10)	$\label{eq:Follow} \mbox{Follow set} \qquad \qquad \mbox{the set of terminals appearing first in any string after}$
	deriving a nonterminal.
	2. Lists
Com	plete 3 of the following lists for full credit. Complete all for extra credit.
(1)	Compiler phases, in order. Briefly describe what each phase does.
	(a) Scanner/Lexer/Tokenizer. Converts a string into a token sequence.
	(b) Parser. Constructs a parse tree.
	(c) Semantic analysis. Type checking.
	(d) Optimizer. Improve performance.
	(e) Code generator. Output code.
(2)	Primitive regular expressions. Briefly describe what each regular expression matches.
	(a) Alternation. a b a or b.
	(b) Catenation. ab a followed by b.
	(c) Kleene Closure a* a 0 or more times.
	(d) Empty Set. {} Nothing.
	(e) Empty String. ""
	(f) Symbol. A character.

NAME:

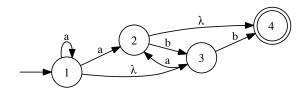
2

(3) Finite automaton elements. Describe each.

	(a) Set of states.
	(b) Start state. The initial state.
	(c) Set of transitions. State x Symbol -> State
	(d) Set of accepting states. (Final states)
	(e) Alphabet. The set of symbols.
(4)	For a grammar to be $LL(1)$, it must be:
	$(\ensuremath{\mathrm{a}})$ Unambiguous, recognized by a deterministic pushdown automaton
	(b) Free of left recursion
	(c) Free of common prefixes
	(d) First(A) disjoint from Follow (A)
	3. FILL IN THE BLANK
Com	
Com	plete the following statements for full credit.
(1)	A pushdown automaton is a finite automaton with one stack.
(2)	A Turing machine is a finite automation with an infinite read/write tape.
(3)	It is never possible to define an NFA which cannot be converted into a DFA.
¹ Left.	right, Leftmost derivation, 1 token lookahead
	• ,

4. Regular languages

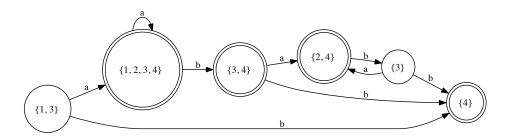
Refer to the Figure below. Answer 3 of the following questions. Answer all for extra credit.



(1) What is the initial state of the DFA using subset construction?

{1,3}

(2) Draw the equivalent DFA using subset construction.



(3) Write the equivalent regular expression.

$$b|aa*(\lambda|b|bb|ba(ba)*(b|\lambda))$$

(4) IPv4 addresses are written as four integers, separated by dots (e.g., 173.203.204.223). Each integer ranges from 0 to 255. Write a regular expression to match precisely these addresses.

$$(25[0-5]|2[0-4][0-9]|1[0-9]\{1,2\}|[0-9] \backslash) \{3\} \\ 25[0-5]|2[0-4][0-9]|1[0-9]\{1,2\}|[0-9] \backslash) \{3\} \\ 25[0-5]|2[0-4][0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[0-9]|1[$$

5

5. Context-free languages

Refer to the context-free grammar below. S is the start symbol. Answer 4 of the following questions. Answer all for extra credit.

$$\begin{split} S \to T & T \to \mathbf{x} \\ S \to S + T & T \to \mathbf{y} \\ S \to S - T & T \to \mathbf{z} \\ S \to S * T & T \to (S) \end{split}$$

(1) Is the grammar above ambiguous? Why or why not?

Unambiguous, because the (symbol forces the parse tree to be formed one way.

(2) Explain why the grammar above is not LL(1).

Left recursion, common prefixes.

(3) What is First(T)?

x, y, z, (

(4) What is Follow(S)?

+, -, *, /,), eof

(5) Perform a leftmost derivation of the following string: $\mathbf{x} * (\mathbf{y} + \mathbf{z})$

6. Extra credit

Complete any of the following for extra credit.

 $Article \rightarrow \mathbf{the}$

(1) List all possible Sentences that can be matched by the grammar below.

 $Sentence \rightarrow NounPhrase \ VerbPhrase \\ NounPhrase \rightarrow Article \ Noun \\ VerbPhrase \rightarrow Verb \ NounPhrase \\ Verb \rightarrow \textbf{kicked} \\$

the boy kicked the ball the boy kicked the boy the ball kicked the boy the ball kicked the ball

(2) Rewrite the grammar on the previous page to be LL(1). Must eliminate left recursion, common prefixes.

S -> T Y

T -> x

-> y

-> z

-> (S)

Y -> + T

-> - T

-> * T

-> / T

-> lambda