

## LECTURE 15

In this lecture, we define and study *stable*  $\infty$ -categories, which are exactly  $\infty$ -categories of the form  $\mathbf{Sptr}(\mathbf{C})$ .

### 1. STABILITY

**Definition 1.1.** Let  $\mathbf{C}$  be a pointed  $\infty$ -category. A **triangle** in  $\mathbf{C}$  is a diagram  $\Delta^1 \times \Delta^1 \rightarrow \mathbf{C}$  depicted as

$$(1.1) \quad \begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow & & \downarrow g \\ 0 & \longrightarrow & Z \end{array}$$

where  $0 \in \mathbf{C}$  is the zero object. We say such a triangle is a **fiber sequence** if it is a pullback square, and a **cofiber sequence** if it is a pushout square.

For a morphism  $f : X \rightarrow Y$ , a **cofiber of  $f$**  is a fiber sequence (1.1). Dually, for a morphism  $g : Y \rightarrow Z$ , a **fiber of  $g$**  is a cofiber sequence (1.1).

1.2. We often abuse notation and write a triangle as  $X \xrightarrow{f} Y \xrightarrow{g} Z$ .

**Warning 1.3.** The datum of a triangle (1.1) is not determined by the chain  $X \xrightarrow{f} Y \xrightarrow{g} Z$ , even up to homotopy. Indeed, knowing such a triangle is equivalent to knowing a null-homotopy of  $g \circ f$ , which is not unique even up to homotopy.

1.4. Note however that a fiber sequence (1.1) is essentially uniquely determined by the morphism  $g$ . Dually, a cofiber sequence (1.1) is essentially uniquely determined by the morphism  $f$ . Hence we can use the notations

$$\mathrm{Fib}(g), \mathrm{Cofib}(f) \in \mathbf{C}$$

as long as we incorporate (1.1) as data in their definitions.

**Definition 1.5.** An  $\infty$ -category  $\mathbf{C}$  is **stable** if it satisfies the following conditions:

- it is pointed
- any morphism in  $\mathbf{C}$  admits a fiber and a cofiber
- a triangle in  $\mathbf{C}$  is a fiber sequence iff it is a cofiber sequence.

1.6. For stable  $\infty$ -categories, we can use the words **fiber-cofiber sequences**.

**Exercise 1.7.** Find all ordinary categories that are stable when viewed as  $\infty$ -categories.

**Proposition 1.8.** Let  $\mathbf{C}$  be a pointed  $\infty$ -category admitting both finite limits and colimits. Then the following conditions are equivalent.

- (i) The functor  $\Sigma : \mathbf{C} \rightarrow \mathbf{C}$  is fully faithful.
- (ii) Any cofiber sequence in  $\mathbf{C}$  of the form  $X \rightarrow 0 \rightarrow Z$  is also a fiber sequence.

- (iii) Any cofiber sequence in  $\mathcal{C}$  is a fiber sequence.
- (iv) Any pushout square in  $\mathcal{C}$  is a pullback square.

**Corollary 1.9.** *Let  $\mathcal{C}$  be a pointed  $\infty$ -category admitting both finite limits and colimits. Then the following conditions are equivalent.*

- (i) The functors  $\Sigma : \mathcal{C} \rightleftarrows \mathcal{C} : \Omega$  are equivalences.
- (ii) A triangle in  $\mathcal{C}$  of the form  $X \rightarrow 0 \rightarrow Z$  is a cofiber sequence iff it is a fiber sequence.
- (iii) A triangle in  $\mathcal{C}$  is a cofiber sequence iff it is a fiber sequence.
- (iv) A square in  $\mathcal{C}$  is a pushout iff it is a pullback.

*Proof of Proposition 1.8.* The implications (ii)  $\Leftarrow$  (iii)  $\Leftarrow$  (iv) are obvious. It remains to show (i)  $\Rightarrow$  (iv). Suppose  $\Sigma$  is fully faithful. Since  $\Omega$  is a right adjoint of  $\Sigma$ , we have  $\text{Id}_{\mathcal{C}} \simeq \Omega \circ \Sigma$ . For a pushout square

$$\begin{array}{ccc} W & \longrightarrow & X \\ \downarrow & & \downarrow \\ Y & \longrightarrow & Z, \end{array}$$

we need to show  $W \rightarrow X \times_Z Y$  is invertible. Consider the following commutative diagram

$$\begin{array}{ccccc} W & \longrightarrow & X & \longrightarrow & 0 \\ \downarrow & & \downarrow & & \downarrow \\ Y & \longrightarrow & Z & \xrightarrow{\quad \cdot \quad} & \Sigma W \\ \downarrow & & & & \downarrow \\ 0 & \longrightarrow & \Sigma W & & \end{array}$$

Here both the inner and outer squares are pushout squares, hence there exists an essentially unique dotted arrow  $Z \rightarrow \Sigma W$  making the above diagram commute. By functoriality of pullbacks, we obtain a morphism  $X \times_Z Y \rightarrow 0 \times_{\Sigma W} 0 \simeq \Omega \Sigma W$  fitting into the following commutative diagram

$$\begin{array}{ccc} W & \longrightarrow & X \times_Z Y \\ \downarrow \simeq & \swarrow \cdot & \downarrow \simeq \\ \Omega \Sigma W & \longrightarrow & \Omega \Sigma X \times_{\Omega \Sigma Z} \Omega \Sigma Y, \end{array}$$

where the vertical morphisms are isomorphisms because of  $\text{Id}_{\mathcal{C}} \simeq \Omega \circ \Sigma$ . By the 2-out-of-6 property of isomorphisms, we obtain  $X \xrightarrow{\simeq} X \times_Z Y$  as desired.  $\square$

**Proposition 1.10.** *Let  $\mathcal{C}$  be a pointed  $\infty$ -category. The following conditions are equivalent.*

- (a) The  $\infty$ -category  $\mathcal{C}$  is stable.
- (b) The  $\infty$ -category  $\mathcal{C}^{\text{op}}$  is stable.
- (c) The  $\infty$ -category  $\mathcal{C}$  admits finite colimits and  $\Sigma : \mathcal{C} \rightarrow \mathcal{C}$  is an equivalence.
- (d) The  $\infty$ -category  $\mathcal{C}$  admits finite limits and  $\Omega : \mathcal{C} \rightarrow \mathcal{C}$  is an equivalence.
- (e) The  $\infty$ -category  $\mathcal{C}$  admits finite colimits and limits, and a square in  $\mathcal{C}$  is a pushout square iff it is a pullback square.

- (f) The  $\infty$ -category  $\mathbf{C}$  admits finite limits and  $\Omega^\infty : \mathbf{Sptr}(\mathbf{C}) \rightarrow \mathbf{C}$  is an equivalence.

*Proof.* The equivalence (a) $\Leftrightarrow$ (b) is obvious. The equivalence (d) $\Leftrightarrow$ (f) was proved last time. It remains to show (c) $\Leftrightarrow$ (a) $\Leftrightarrow$ (e) because (d) $\Leftrightarrow$ (b) would follow by passing to the opposite  $\infty$ -category.

Suppose  $\mathbf{C}$  admits finite colimits and limits, then (c) $\Leftrightarrow$ (a) $\Leftrightarrow$ (e) follow from (i) $\Leftrightarrow$ (iii) $\Leftrightarrow$ (iv) in Proposition 1.8 (and its dual version). Hence it remains to show (a) or (c) implies  $\mathbf{C}$  admits finite colimits and limits.

For (a), we only need to show a stable  $\infty$ -category admits coequalizers. This follows from Exercise 2.13 below.

For (c), we can use  $\iota : \mathbf{C} \rightarrow \mathbf{Ind}(\mathbf{C})$  to embed  $\mathbf{C}$  into a presentable  $\infty$ -category. Note that  $\mathbf{Ind}(\mathbf{C})$  is pointed because  $\iota$  preserves and detects both finite colimits and limits. Moreover,  $\Sigma_{\mathbf{Ind}(\mathbf{C})}$  can be identified with  $\mathbf{Ind}(\Sigma_{\mathbf{C}}) : \mathbf{Ind}(\mathbf{C}) \rightarrow \mathbf{Ind}(\mathbf{C})$  and thereby is also an equivalence. It follows from the previous discussion that  $\mathbf{Ind}(\mathbf{C})$  satisfies all the properties in the proposition. In particular, any pushout square in  $\mathbf{Ind}(\mathbf{C})$  is a pullback square. Since  $\iota$  preserves and detects both finite colimits and limits, we see the same holds for  $\mathbf{C}$ . In particular, it admits pullbacks and therefore all finite limits as desired.  $\square$

**Corollary 1.11.** *Let  $\mathbf{C}$  be a pointed  $\infty$ -category that admits finite limits. The  $\infty$ -category  $\mathbf{Sptr}(\mathbf{C})$  is stable.*

**Exercise 1.12.** *Let  $\mathbf{C}$  be a stable  $\infty$ -category. Then  $f : X \rightarrow Y$  is an isomorphism iff  $\mathrm{Fib}(f)$  is a zero object iff  $\mathrm{Cofib}(f)$  is a zero object.*

**Exercise 1.13.** *Let  $f : X \rightarrow Y$  be a morphism in a stable  $\infty$ -category  $\mathbf{C}$ . Show that  $\Sigma \mathrm{Fib}(f) \simeq \mathrm{Cofib}(f)$ .*

## 2. HOMOTOPY CATEGORY OF STABLE $\infty$ -CATEGORY

2.1. In this section, let  $\mathbf{C}$  be a stable  $\infty$ -category.

**Exercise 2.2.** *Show that the canonical morphism  $X \sqcup Y \rightarrow X \times Y$  is invertible. Hint:*

$$\begin{array}{ccc} X \sqcup Y & \longrightarrow & X \sqcup 0 \\ \downarrow & & \downarrow \\ 0 \sqcup Y & \longrightarrow & 0 \sqcup 0 \end{array}$$

*is a pushout square.*

2.3. Since there is a canonical equivalence between  $X \sqcup Y$  and  $X \times Y$ , we use  $X \oplus Y$  to denote both of them.

**Exercise 2.4.** *Let  $f, g : X \rightarrow Y$  be two morphisms. Show that the composition*

$$X \rightarrow X \oplus X \xrightarrow{(f,g)} Y \oplus Y \rightarrow Y$$

*gives a well-defined binary operator on  $\pi_0 \mathrm{Hom}_{\mathbf{C}}(X, Y)$ . We denote the above composition by  $f + g$ .*

**Exercise 2.5.** *Let  $f, g : X \rightarrow Y$  be two morphisms. Show that the above binary operator coincides with the addition operator on the abelian group*

$$\pi_0 \mathrm{Hom}_{\mathbf{C}}(X, Y) \simeq \pi_0 \mathrm{Hom}_{\mathbf{C}}(X, \Omega^2 \Sigma^2 Y) \simeq \pi_0 \Omega^2 \mathrm{Hom}_{\mathbf{C}}(X, \Sigma^2 Y) \simeq \pi_2 \mathrm{Hom}_{\mathbf{C}}(X, \Sigma^2 Y).$$

**Exercise 2.6.** Let  $\sigma : \Delta^1 \times \Delta^1 \rightarrow \mathbf{C}$  be a diagram of the form

$$\begin{array}{ccc} X & \longrightarrow & 0 \\ \downarrow & & \downarrow \\ 0 & \longrightarrow & \Sigma Y \end{array}$$

and  $\sigma'$  be its transpose. By the universal property of pushouts,  $\sigma$  induces a morphism  $\eta_\sigma : X \rightarrow \Omega \Sigma Y \simeq Y$ , which is an well-defined element in  $\pi_0 \text{Hom}_{\mathbf{C}}(X, Y)$ . Show that  $\eta_\sigma + \eta_{\sigma'} = 0$ .

**Corollary 2.7.** The homotopy category  $\mathbf{hC}$  is an additive category.

2.8. From now on, we write  $X[n] := \Sigma^n X$ , where for  $n < 0$  we take  $\Sigma^n := \Omega^{-n}$ . Note that these objects are well-defined up to homotopy.

**Definition 2.9.** Let  $X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} X[1]$  be a chain in  $\mathbf{hC}$ . We say it is a **distinguished triangle in  $\mathbf{hC}$**  if there exists a diagram  $\Delta^1 \times \Delta^2 \rightarrow \mathbf{C}$  depicted as

$$\begin{array}{ccccc} X & \xrightarrow{\tilde{f}} & Y & \longrightarrow & 0 \\ \downarrow & & \downarrow \tilde{g} & & \downarrow \\ 0' & \longrightarrow & Z & \longrightarrow & W \end{array}$$

such that

- the objects  $0$  and  $0'$  are zero objects
- the morphisms  $\tilde{f}$  and  $\tilde{g}$  lift  $f$  and  $g$  respectively
- the outer square is a fiber-cofiber sequence
- the composition  $Z \rightarrow W \xrightarrow{\cong} X[1]$  (which is well-defined up to homotopy) lifts  $h$ .

**Theorem 2.10** (HA.1.1.2.14). The above choice of the translation functor and the distinguished triangles makes  $\mathbf{hC}$  a triangulated category.

**Exercise 2.11.** What would happen if we use diagrams of the form

$$\begin{array}{ccc} X & \longrightarrow & 0 \\ \downarrow & & \downarrow \\ Y & \longrightarrow & Z \\ \downarrow & & \downarrow \\ 0' & \longrightarrow & W \end{array}$$

to define distinguished triangles in  $\mathbf{hC}$ ?

**Exercise 2.12.** For  $X \in \mathbf{C}$ , construct a fiber-cofiber sequence

$$X \xrightarrow{(\text{id}, -\text{id})} X \oplus X \xrightarrow{(\text{id}, \text{id})} X.$$

**Exercise 2.13.** Show that the coequalizer of  $f, g : X \rightrightarrows Y$  is canonically equivalent to  $\text{Cofib}(f - g)$ . *Hint:*

$$\begin{array}{ccccc} X & \xrightarrow{(\text{id}, -\text{id})} & X \oplus X & \xrightarrow{(f, g)} & Y \\ \downarrow & & \downarrow (\text{id}, \text{id}) & & \downarrow \\ 0 & \longrightarrow & X & \longrightarrow & \bullet \end{array}$$

### 3. MAPPING SPECTRA

**Construction 3.1.** Let  $\mathcal{C}$  be a stable  $\infty$ -category. For  $X, Y \in \mathcal{C}$ , define

$$\underline{\text{Hom}}(X, Y)_n := \text{Maps}(X, Y[n]).$$

Note that for any  $n \geq 0$ , we have an isomorphism

$$\alpha_n : \text{Maps}(X, Y[n]) \simeq \text{Maps}(X, \Omega Y[n+1]) \simeq \Omega \text{Maps}(X, Y[n+1]).$$

Let  $\underline{\text{Hom}}(X, Y) \in \mathbf{Sptr}$  be the spectrum given by the spaces  $\{\underline{\text{Hom}}(X, Y)_n\}$  and the isomorphisms  $\alpha_n$ . We call it the **mapping spectrum** between  $X$  and  $Y$ .

**Remark 3.2.** The above spectrum  $\underline{\text{Hom}}(X, Y)$  is well-defined up to homotopy. In future lectures, we will equip  $\mathbf{Sptr}$  with a symmetric monoidal structure such that any stable  $\infty$ -category  $\mathcal{C}$  is canonically enriched over  $\mathbf{Sptr}$ .

**Definition 3.3.** Let  $\mathcal{C}$  be a stable  $\infty$ -category. For  $X, Y \in \mathcal{C}$ , define

$$\text{Ext}^n(X, Y) := \pi_0 \underline{\text{Hom}}(X, Y)_n \simeq \pi_0 \text{Maps}(X, Y[n])$$

and call it the  $n$ -th extension group between  $X$  and  $Y$ .

3.4. Note that the extension groups only depend on the images of  $X$  and  $Y$  in the triangulated category  $\mathbf{hC}$ . It is well-known that for any distinguished triangle  $Y_0 \rightarrow Y_1 \rightarrow Y_2 \rightarrow Y_0[1]$  in  $\mathbf{hC}$ , we have a long exact sequence

$$\cdots \rightarrow \text{Ext}^n(X, Y_0) \rightarrow \text{Ext}^n(X, Y_1) \rightarrow \text{Ext}^n(X, Y_2) \rightarrow \text{Ext}^{n+1}(X, Y_0) \rightarrow \cdots$$

In the next lecture, we will construct these long exact sequences using the theory of stable  $\infty$ -categories.

### 4. EXACT FUNCTORS

4.1. Proposition 1.10 and Exercise 2.13 imply the following result.

**Proposition-Definition 4.2.** Let  $F : \mathcal{C} \rightarrow \mathcal{C}'$  be a functor between stable  $\infty$ -category. The following conditions are equivalent.

- The functor  $F$  preserves zero objects and fiber-cofiber sequence.
- The functor  $F$  is left exact, i.e., preserves finite limits.
- The functor  $F$  is right exact, i.e., preserves finite colimits.

We say  $F$  is **exact** if it satisfies the above conditions.

**Exercise 4.3.** Let  $F : \mathcal{C} \rightarrow \mathcal{C}'$  be an exact functor between stable  $\infty$ -category. Show that  $\mathbf{h}F : \mathbf{hC} \rightarrow \mathbf{hC}'$  has a natural structure of an exact functor<sup>1</sup> between triangulated categories.

<sup>1</sup>Also known as a triangulated functor. Warning: being a triangulated functor is a structure rather than property.

**Definition 4.4.** Let  $\text{Cat}_\infty^{\text{ex}} \subseteq \text{Cat}_\infty$  be the sub- $\infty$ -category of small stable  $\infty$ -categories and exact functors between them.

**Exercise 4.5.** Let  $\mathcal{C}$  be a stable  $\infty$ -category. Show that  $\Sigma$  and  $\Omega$  are exact. What are the triangulated functors induced by them?

## 5. CLOSURE PROPERTIES

**Exercise 5.1.** Let  $\mathcal{C}$  be a stable  $\infty$ -category and  $K$  be a simplicial set. Show that  $\text{Fun}(K, \mathcal{C})$  is stable.

**Exercise 5.2.** Let  $\mathcal{C}$  and  $\mathcal{D}$  be stable  $\infty$ -categories. Show that  $\text{Fun}_{\text{ex}}(\mathcal{C}, \mathcal{D})$  is stable.

**Exercise 5.3.** Let  $\mathcal{C}$  and  $\mathcal{D}$  be presentable  $\infty$ -categories such that  $\mathcal{D}$  is stable. Show that  $\text{LFun}(\mathcal{C}, \mathcal{D})$  is stable.

**Exercise 5.4.** Let  $\mathcal{C}$  be a small stable  $\infty$ -category, then  $\text{Ind}(\mathcal{C})$  is stable.

**Exercise 5.5.** Let  $\mathcal{C}$  be a stable  $\infty$ -category, then  $\mathcal{C}^{\text{cpt}}$  is stable. In particular,  $\text{Sptr}^{\text{fin}}$  is stable.

**Exercise 5.6.** Let  $\mathcal{C}$  be a stable  $\infty$ -category, show that the idempotent completion of  $\mathcal{C}$  is also stable.

**Exercise 5.7.** Let  $\mathcal{C}$  be a stable  $\infty$ -category, is  $\text{PShv}(\mathcal{C})$  stable?

**Theorem 5.8.** The  $\infty$ -category  $\text{Cat}_\infty^{\text{ex}}$  admits small limits and the inclusion  $\text{Cat}_\infty^{\text{ex}} \rightarrow \text{Cat}_\infty$  preserves and detects small limits.

*Sketch.* We only need to show for any small diagram  $K \rightarrow \text{Cat}_\infty$ ,  $i \mapsto \mathcal{C}_i$  such that each  $\mathcal{C}_i$  is stable and each connecting functor  $\mathcal{C}_i \rightarrow \mathcal{C}_j$  is exact, we have

- the limit  $\infty$ -category  $\mathcal{C} := \lim_i \mathcal{C}_i$  is stable
- the evaluating functors  $\mathcal{C} \rightarrow \mathcal{C}_i$  are exact.

Both claims can be checked using the explicit description of objects and mapping spaces in  $\mathcal{C}$ .  $\square$

5.9. Similarly, one can prove the following result.

**Theorem 5.10.** The  $\infty$ -category  $\text{Cat}_\infty^{\text{ex}}$  admits small filtered colimits and the inclusion  $\text{Cat}_\infty^{\text{ex}} \rightarrow \text{Cat}_\infty$  preserves and detects small filtered colimits.

## 6. A UNIVERSAL PROPERTY OF $\text{Sptr}$

6.1. The following result implies  $\text{Sptr}$  is the stable  $\infty$ -category freely generated by one object under small colimits.

**Exercise 6.2.** Let  $\mathcal{D}$  be a presentable stable  $\infty$ -category. Show that evaluating at  $\mathbb{S} \in \text{Sptr}$  induces an equivalence

$$\text{LFun}(\text{Sptr}, \mathcal{D}) \xrightarrow{\sim} \mathcal{D}.$$

*Hint:* show  $\text{RFun}(\mathcal{D}, \text{Sptr}) \xrightarrow{\Omega^\infty \circ -} \text{RFun}(\mathcal{D}, \text{Spc})$  is an equivalence.

**Exercise 6.3.** Let  $\mathcal{D}$  be a presentable stable  $\infty$ -category. Show that evaluating at  $\mathbb{S} \in \text{Sptr}^{\text{fin}}$  induces an equivalence

$$\text{Fun}_{\text{ex}}(\text{Sptr}^{\text{fin}}, \mathcal{D}) \xrightarrow{\sim} \mathcal{D}.$$

## APPENDIX A. TRIANGULATED CATEGORIES WITHOUT MODELS

- A.1. There are triangulated categories that are not the homotopy category of any stable  $\infty$ -category.
- A.2. There are exact functors between homotopy categories of stable  $\infty$ -categories that do not come from exact functors between the stable  $\infty$ -categories.
- A.3. **Suggested readings.** [MSS07].

## REFERENCES

- [MSS07] Fernando Muro, Stefan Schwede, and Neil Strickland. Triangulated categories without models. *Inventiones mathematicae*, 170(2):231–241, 2007.