

Cryptology Exercise Week 6

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RSA decryption works on the entire domain

Because n is the product of two co-prime numbers p and q , the Chinese Remainder Theorem applies here. Because the f -function in the Chinese Remainder Theorem is injective, in order to show $x^{ed} \equiv x \pmod n$, it suffices to show that $x^{ed} \equiv x \pmod p$ and $x^{ed} \equiv x \pmod q$, for all $x \in \mathbb{Z}_n$.

Here, we are showing the case where $x \notin \mathbb{Z}_n^*$. Since n is the product of the prime numbers p and q , we have that x is either the product of p and q , or a multiple of p or q .

In the case where $x = p \cdot q$, we have that $x \equiv x^{ed} \equiv 0 \pmod p$ and $x \equiv x^{ed} \equiv 0 \pmod q$.

We then discuss the case where x is a multiple of p or q . Without loss of generality, we assume that x is a multiple of p , but not of q . It is obvious that $x \equiv x^{ed} \equiv 0 \pmod p$. Because q is a prime number and x is not a multiple of q , we have that x and q are co-prime and $(x \bmod q) \in \mathbb{Z}_q^*$. Notice that $|\mathbb{Z}_q^*| = q - 1$, hence we have

$$(x \bmod q)^{ed} \equiv x^{ed} \equiv x^{ed \bmod (q-1)} \pmod q$$

Since $ed \equiv 1 \pmod{(p-1)(q-1)}$, we have that $ed \equiv 1 \pmod{(q-1)}$, hence

$$x^{ed} \equiv x \pmod q$$