Cryptology Exercise Week 5

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Exercise 6.6

According to the definition of AES, the first column of R is

$$
\begin{aligned}
& R[:,0] \\
=& MC([S(a_{0,0}), S(a_{1,1}), S(a_{2,2}), S(a_{3,3})]) \\
=& MC([S(a_{0,0}), 0, 0, 0] \\
& \quad \oplus [0, S(a_{1,1}), 0, 0] \\
& \quad \oplus [0, 0, S(a_{2,2}), 0] \\
& \quad \oplus [0, 0, 0, S(a_{3,3})]) \\
=& MC([S(a_{0,0}), 0, 0, 0]) \\
& \oplus MC([0, S(a_{1,1}), 0, 0]) \\
& \oplus MC([0, 0, S(a_{2,2}), 0]) \\
& \oplus MC([0, 0, 0, S(a_{3,3})]) \\
=& T_0(a_{0,0}) \oplus T_1(a_{1,1}) \oplus T_2(a_{2,2}) \oplus T_3(a_{3,3})
\end{aligned}
$$

Similarly, the rest of each column are

$$
\begin{aligned}
R[:,1] &= T_0(a_{0,1}) \oplus T_1(a_{1,2}) \oplus T_2(a_{2,3}) \oplus T_3(a_{3,0}) \\
R[:,2] &= T_0(a_{0,2}) \oplus T_1(a_{1,3}) \oplus T_2(a_{2,0}) \oplus T_3(a_{3,1}) \\
R[:,3] &= T_0(a_{0,3}) \oplus T_1(a_{1,0}) \oplus T_2(a_{2,1}) \oplus T_3(a_{3,2})
\end{aligned}
$$

Implementation

For each of T_0, T_1, T_2, T_3, we create a table that has 256 entries, mapping from the input byte, i.e a_{ij}, to the output 4-byte column. Then in each round of AES, we can simply do table lookups for each byte followed by XOR operations to obtain R and then XOR R with the round key to get the final rsesult of each round. The memory usage of the tables is $4 * 4 * 256\text{B} = 4\text{KB}$.