

Program Logics Hand-in 1

Zijun Yu 202203581

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Exercise 1

Part 1

```
([], [0 ↦ singleton1])  
→ ( [], [0 ↦ inr(ref(1, inl()))])  
→ ([l1 ↦ (1, inl())], [0 ↦ inrl1])
```

Part 2

```
([], [0 ↦ inc(singleton1)])  
→ ([l1 ↦ (1, inl())], [0 ↦ inc(inrl1)])  
→ ([l1 ↦ (1, inl())], [0 ↦  
  match inrl1 with inlx1 ⇒ () | inrx2 ⇒ let v := π1(!x2) in let t := π2(!x2) in x2 ← (v + 1, t); inct end])  
→ ([l1 ↦ (1, inl())], [0 ↦ let v := π1(!l1) in let t := π2(!l1) in l1 ← (v + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ let v := π1(1, inl()) in let t := π2(!l1) in l1 ← (v + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ let v := 1 in let t := π2(!l1) in l1 ← (v + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ let t := π2(!l1) in l1 ← (1 + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ let t := π2(1, inl()) in l1 ← (1 + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ let t := inl() in l1 ← (1 + 1, t); inct])  
→ ([l1 ↦ (1, inl())], [0 ↦ l1 ← (1 + 1, inl()); inc inl()])  
→ ([l1 ↦ (1, inl())], [0 ↦ l1 ← (2, inl()); inc inl()])  
→ ([l1 ↦ (2, inl())], [0 ↦ (); inc inl()])  
→ ([l1 ↦ (2, inl())], [0 ↦ inc inl()])  
→ ([l1 ↦ (2, inl())], [0 ↦  
  match inl() with inlx1 ⇒ () | inrx2 ⇒ let v := π1(!x2) in let t := π2(!x2) in x2 ← (v + 1, t); inct end])  
→ ([l1 ↦ (2, inl())], [0 ↦ ()])
```

The syntactic sugars are not expanded in the reduction.

Exercise 2

Approach 1

$$\frac{\overline{P * Q \vdash P} \quad \overline{P * Q \vdash Q}}{P * Q \vdash P \wedge Q}$$

Approach 2

Using the result from Exercise 3.

$$\frac{\overline{P \wedge Q \vdash P \wedge Q} \quad \frac{P \vdash Q \multimap (P \wedge Q) \quad \overline{Q \vdash Q}}{P * Q \vdash P \wedge Q}}{P * Q \vdash P \wedge Q}$$

Exercise 3

$$\frac{P \vdash Q \multimap R \quad \overline{Q \vdash Q}}{P * Q \vdash R}$$

Exercise 4

From left to right:

$$\frac{\frac{\frac{\overline{\Gamma \mid \exists x.P * Q \vdash \exists x.P * Q}}{\Gamma, x \mid P * Q \vdash \exists x.P * Q}}{\Gamma, x \mid Q \vdash P \multimap \exists x.P * Q}}{\Gamma \mid \exists x.Q \vdash P \multimap \exists x.P * Q}}{\Gamma \mid P * \exists x.Q \vdash \exists x.P * Q}$$

From right to left:

$$\frac{\frac{\overline{\Gamma, x \mid P \vdash P} \quad \frac{\overline{\Gamma, x \mid \exists x.Q \vdash \exists x.Q}}{\Gamma, x \mid Q \vdash \exists x.Q}}{\Gamma, x \mid P * Q \vdash P * \exists x.Q}}{\Gamma \mid \exists x.P * Q \vdash P * \exists x.Q}$$