

# Program Logics Hand-in 4

Zijun Yu 202203581

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## Exercise 1

We introduce  $m$  and apply HT-INV-OPEN and E- $\triangleright$ , since  $l \leftarrow (m+1)$  is an atomic expression (assuming that  $m+1$  is a value), SFTS

$$\boxed{I \ n}^\ell \vdash \{\exists m'. \triangleright(m' \geq n \wedge l \hookrightarrow m') * (m \geq n)\} l \leftarrow (m+1) \{ \neg \triangleright(\exists m. m \geq n \wedge l \hookrightarrow m) * \text{True} \}_{\mathcal{E} \setminus \ell} \quad (1)$$

We then use HT-EXIST and introduce  $m'$ , SFTS

$$\boxed{I \ n}^\ell \vdash \{\triangleright(m' \geq n \wedge l \hookrightarrow m') * (m \geq n)\} l \leftarrow (m+1) \{ \neg \triangleright(\exists m. m \geq n \wedge l \hookrightarrow m) * \text{True} \}_{\mathcal{E} \setminus \ell} \quad (2)$$

Using HT-FRAME-ATMOIC with HT-STORE and other structure rules we have

$$\boxed{I \ n}^\ell \vdash \{\triangleright(l \hookrightarrow m' \wedge m \geq n)\} l \leftarrow (m+1) \{v.v = () \wedge l \hookrightarrow (m+1) \wedge m \geq n\}_{\mathcal{E} \setminus \ell} \quad (3)$$

Because the precondition in (2) implies the precondition in (3), we apply HT-CSQ, and we are left to show

$$v = () \wedge l \hookrightarrow (m+1) \wedge m \geq n \vdash \triangleright(\exists m. m \geq n \wedge l \hookrightarrow m) * \text{True}$$

This holds as we can take  $\exists m$  to be  $m+1$  and apply E- $\triangleright$  and LATER-WEAK.

## Exercise 2

$$\frac{\frac{P_2 \vdash \models Q}{P_1 * P_2 \vdash P_1 * \models Q \vdash \models (P_1 * Q)} \quad \frac{P_1 * Q \vdash \models R}{\models (P_1 * Q) \vdash \models \models R \vdash \models R}}{P_1 * P_2 \vdash \models R}$$

## Exercise 3

### HT-TOKEN-ALLOC

By HT-CSQ-VS, it suffices to show  $S \vdash P \Rightarrow \exists \gamma. \overset{\neg}{\underset{\neg}{\text{T}}}^\gamma * P$ , which is equivalent to

$$S \vdash \Box(P \Rightarrow \models (\exists \gamma. \overset{\neg}{\underset{\neg}{\text{T}}}^\gamma * P))$$

Because  $S \vdash \text{True} \vdash \Box \text{True}$ , by TRANS, it suffices to show

$$\Box \text{True} \vdash \Box(P \Rightarrow \models (\exists \gamma. \overset{\neg}{\underset{\neg}{\text{T}}}^\gamma * P))$$

Then by PERSISTENTLY-INTRO, it suffices to show

$$\Box \text{True} \vdash P \Rightarrow \models (\exists \gamma. \overset{\neg}{\underset{\neg}{\text{T}}}^\gamma * P)$$

which simplifies to

$$P \vdash \models (\exists \gamma. [\bar{T}]^{\gamma} * P)$$

which can be derived from GHOST-ALLOC, \*I, and UPD-FRAME.

### HT-TOKEN-UPDATE-PRE

By HT-CSQ-VS, it suffices to show

$$S \vdash [\bar{S}]^{\gamma} * P \Rightarrow [\bar{F}]^{\gamma} * P$$

By applying the same trick as in HT-TOKEN-ALLOC, it suffices to show

$$[\bar{S}]^{\gamma} * P \vdash \models ([\bar{F}]^{\gamma} * P)$$

which follows from GHOST-UPDATE, \*I, and UPD-FRAME.

### HT-TOKEN-UPDATE-POST

By HT-CSQ-VS, what we are left to show is exactly the same as in HT-TOKEN-UPDATE-PRE.

## Exercise 4

By HT-LOAD, we know that

$$\{\triangleright l \hookrightarrow n\} ! l \{v.v = n * l \hookrightarrow n\}$$

By HT-CSQ, we further have

$$\{\triangleright l \hookrightarrow n\} ! l \left\{ v. \left( v = n \vee v = n + 1 * [\bar{F}]^{\gamma_1} \right) * l \hookrightarrow n \right\}$$

Then by HT-FRAME-ATOMIC and \*I, we know that

$$\left\{ \triangleright \left( l \hookrightarrow n * [\bar{S}]^{\gamma_1} \right) * [\bar{\frac{1}{2}}]^{\gamma_2} \right\} ! l \left\{ v. \left( v = n \vee v = n + 1 * [\bar{F}]^{\gamma_1} \right) * l \hookrightarrow n * [\bar{S}]^{\gamma_1} * [\bar{\frac{1}{2}}]^{\gamma_2} \right\}$$

Because  $l \hookrightarrow n * [\bar{S}]^{\gamma_1} \vdash \triangleright (l \mapsto n * [\bar{S}]^{\gamma_1}) \vdash \triangleright I(\gamma_1, \gamma_2, n)$ , by HT-CSQ, we have that

$$\left\{ \triangleright \left( l \hookrightarrow n * [\bar{S}]^{\gamma_1} \right) * [\bar{\frac{1}{2}}]^{\gamma_2} \right\} ! l \left\{ v. \left( v = n \vee v = n + 1 * [\bar{F}]^{\gamma_1} \right) * \triangleright I(\gamma_1, \gamma_2, n) * [\bar{\frac{1}{2}}]^{\gamma_2} \right\}$$

which is what we need to show.