

RANDORI SEC

Using n-day to break kirin's
secureboot for fun and profit

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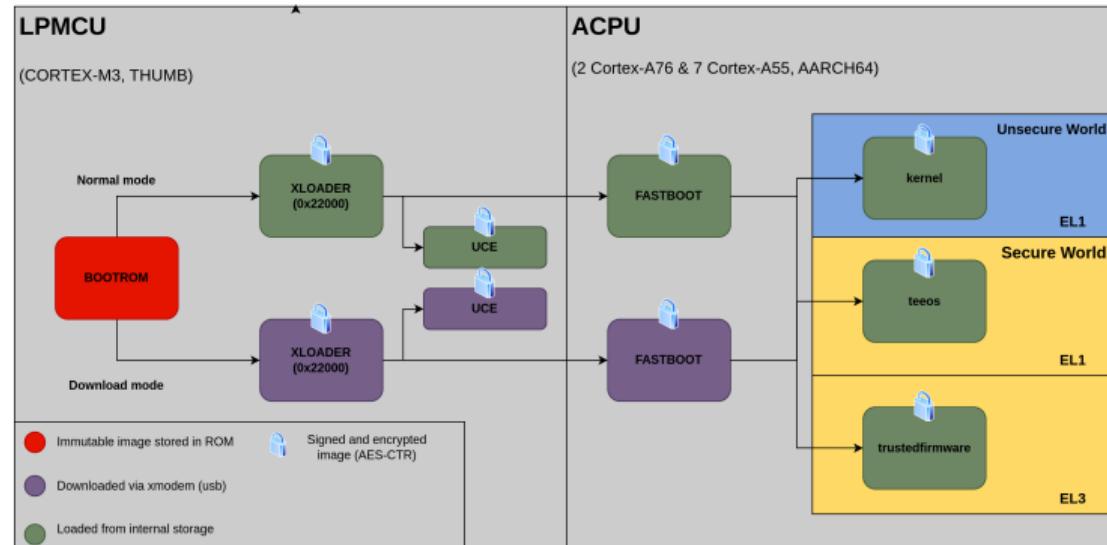


Target device



- ▶ Device: Huawei P40 Lite.
- ▶ SoC: kirin 810

Kirin's bootchain



Vendor's protocol: xmodem

HEADFRAME

CMD (0xFE) 1 bytes	Sequence 1 byte	Sequence~ 1 byte	filetype 1 byte	length 4 bytes (uint)	address 4 bytes (uint)	CRC 2 bytes (ushort)
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DATAFRAME

CMD (0xDA) 1 bytes	Sequence 1 byte	Sequence~ 1 byte	Data up to 1024 bytes	CRC 2 bytes (ushort)
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TAILFRAME

CMD (0xED) 1 bytes	Sequence 1 byte	Sequence~ 1 byte	CRC 2 bytes (ushort)
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Response

ERR	Invalid frame field value
NAK	CRC check failed
ACK	Everything is good

INQUIRY FRAME

CMD (0xCD) 1 bytes	Sequence 1 byte	Sequence~ 1 byte	CRC 2 bytes (ushort)
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Summary

- Vulnerability root cause analysis.
- Turning arbitrary write into arbitrary execution
- Decrypting bootchain's images.
 - Setting up decryption oracle.
 - Patching ROM instructions with FPB.
- Patching images to bypass secureboot.

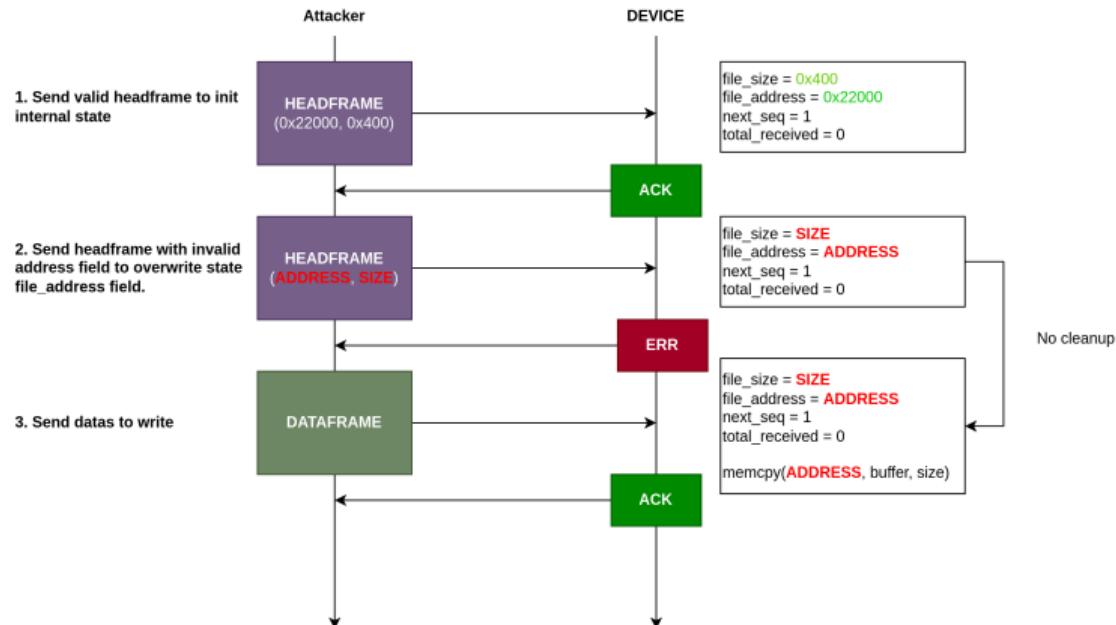
CVE-2021-22433: Head Re-Send State Machine Configuration

```
// Fill state struct
file_addr = _byteswap_ulong(*(_DWORD *)&usb->ss_bulk_buf[8]);
file_size = _byteswap_ulong(*(_DWORD *)&usb->ss_bulk_buf[4]);
usb->file_type = file_type;
usb->file_size = file_size;
usb->file_address = (_BYTE *)file_addr; ←
usb->file_complete = file_addr;
// Check download address
if ( file_addr == 0x22000 ) ←
{
    if ( (file_size & 0x3FF) != 0 )
        size = 2;
    else
        size = 1;
    usb->total_received = 0;
    usb->latest_seen_seq = 0;
    usb->total_frame_count = size + (file_size >> 10);
    usb->next_seq = 1;
}
else
{
    printf("[USBE]file load addr error:%x\n", file_addr); ←
    curr_frame_seq = 7;
}
```

Credits to TASZK Security Labs for discovery and exploitation

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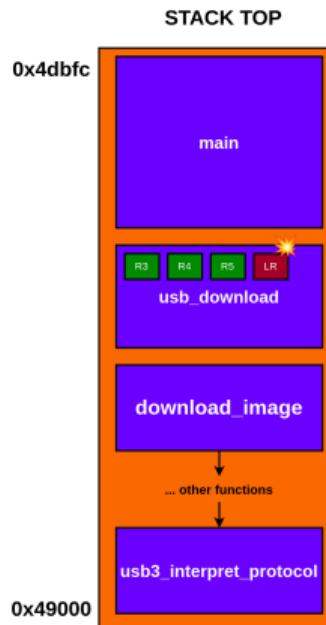
CVE-2021-22433: Triggering vulnerability



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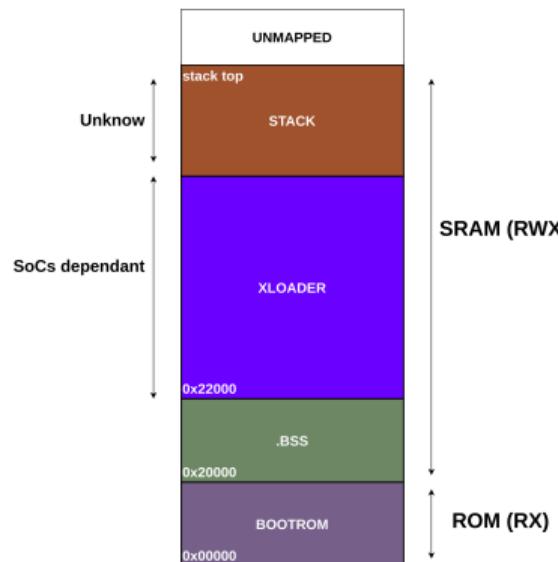
Gaining arbitrary code execution



- ▶ Highly deterministic!
 - ▶ Offset from stack top is easily predictable.
 - ▶ Call happens early in bootrom.
 - ▶ We can smash first 52 bytes of the stack to overwrite saved PC.
- ▶ Same technique can be used on other SoCs.

Gaining arbitrary code execution: where is stack top ?

Hopefully we can reduce range of values to bruteforce



Model	SoC	Xloader size	Stack top
POT	kirin 710	0x2d000	0x49bfcc
JNY	kirin 810	0x28000	?
EML	kirin 970	0x22000	0x4dbfc
YAL	kirin 980	0x27000	0x4dbfc

Gaining arbitrary code execution: finding usb_download with shellcode

```
4F F4 08 30      MOV.W   R0, #0x22000 ; image_addr
FF F7 35 FE      BL      usb_download
                  ; CODE XREF: start+142:p
usb_download      PUSH    {R3-R5,LR}    ; Save registers on stack
38 B5             MOV     R5, R0
05 46             BL      flip_bootrom_flag
04 F0 C2 FD       BL      sub_4D90
04 F0 94 FC       BL      flip_bootrom_flag
04 F0 BE FD       MOVS   R0, #0      ; result
00 20             BL      sub_4FCC
04 F0 AD FD       BL      flip_bootrom_flag
04 F0 B9 FD       BL      download_image
03 F0 7F F9
```

- ▶ **usb_download** open/close serial connection.
- ▶ Search a pattern in bootrom to find the function.

Gaining arbitrary code execution

Memory at 0x4ebfc is not mapped (cause device to crash), which mean that stack top is at 0x4dbfc

```
1 [+] Trying to write at address 0x4cbfc
2 [+] Trying to write at address 0x4dbfc
3 [+] Trying to write at address 0x4ebfc
4 [!] Failed to send dataframe 1
5 [!] Failed to write at address 0x4ebfc
6 Press a key to continue ...
```

Our shellcode found **usb_download** at 0x045C. By using inquiry frame, we sucessfully dumped bootrom in 4h.

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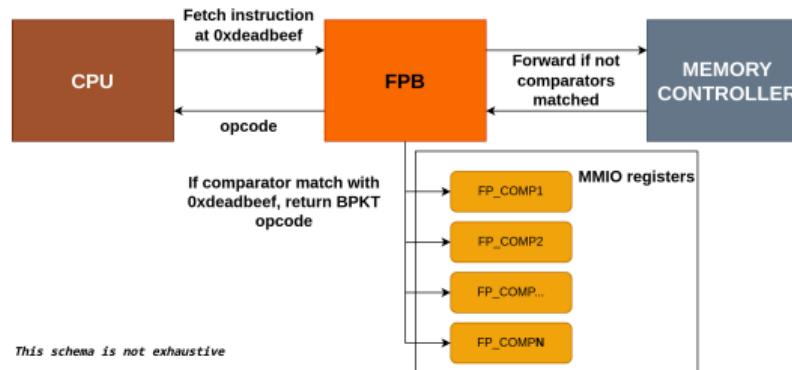
Decrypting images

We need to call these 2 functions :

- ▶ **AES_Init** located at 0x00006968. Setup cryptocell in AES-CTR mode and fill nonce + counter.
- ▶ **AES_Update** located at 0x000069E4. Decrypt buffer of arbitrary size using AES-CTR.

But we can only read 4 bytes per 4 bytes which reduced decryption ...

Increasing decryption speed: Hooking with FPB



- ▶ Patch memory access in ROM and SRAM.
 - ▶ Insert breakpoints in bootrom.
- ▶ Hijack execution flow and change registers content.

Increasing decryption speed: Hooking with FPB

```
inquiry_val = get_inquiry_val();
xmodem_send_inquiry(usb, inquiry_val);
return 0;

int __fastcall xmodem_send_inquiry(struct usb3_pcd *usb, int inquiry)
{
    usb->ss_bulk_buf[0] = inquiry; ←
    usb->ss_bulk_buf[2] = BYTE2(inquiry);
    usb->ss_bulk_buf[3] = HIBYTE(inquiry);
    usb->in_ep.req.length = 4; ←
    usb->in_ep.req.bufdma = usb->ss_bulk_buf;
    usb->ss_bulk_buf[1] = BYTE1(inquiry);
    usb->in_ep.req.complete = (void (__fastcall *)(void *))ep_req_complete;
    return usb3_bulk_out_continue_transfer((int)usb, (int)&usb->in_ep);
}
```

- ▶ Hook **xmodem_send_inquiry** using FPB and breakpoints
- ▶ Inside debug handler we have to :
 - ▶ Change size of usb data to be sent to host
 - ▶ Change content of usb data to be sent to host

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Secureboot bypass: Bootrom stage

```
46     usb_err = usb_download(0x22000u);
47     if ( usb_err )
48         break;
49 verify_xloader_signature:
50     if ( al || (init_err = init_cc()) == 0 )
51     {
52         counter = 0x50C7A5C3;
53         for ( i = 0; i != 2; ++i )
54         {
55             verify_err = cc_img_verification(0x22000, i);
56             if ( verify_err )
57             {
58                 if ( verify_err != 0x50C7B1E4 )
59                     update_inquiry_by_err(verify_err);
60                 al = 1;
61                 goto LABEL_18;
62             }
63             j = 3;
64             do
65             {
66                 ++counter;
67                 udelay_100us(1);
68                 --j;
69             }
70             while ( j );
71         }
72         if ( counter != 0x50C7A5C9 )
73         {
74             update_inquiry(0x50);
75             al = 1;
76             goto LABEL_18;
77         }
78         goto jmp_xloader; ←
79     }
```

- ▶ Patch call to cc_image_verification.
- ▶ Bootrom will naturally go to label **jmp_xloader**.

Secureboot bypass: Xloader stage

```

printf("xloader main download mode to read uce!\n");
v4 = (v2 + 52 + 32 * phdr_idx);           // parse ELF header
expected_address = v4->p_paddr;
if ( expected_address != download_image(0) )
{
    printf("UCEStgErr0x%x\n", v4->p_paddr);
    system_reset(162, 1);
}
if ( UceDlVerify(phdr_idx, lcs) != -1 )
    goto done;
printf("UceDlVerify fail!\n");

```

load_uce function

```

err = sub_33050(&v10, FASTBOOT, 0x5FE000);
if ( err != -1 )
{
    result = err + 15;
    if ( !result )
        return result;
    if ( memcmp_s(0xBFC20000, 4096, 0x1C000000, 0x1000u) )
    {
        printf("vrl cpy E\n");
    }
    else if ( sub_23A02(0x1C000000, 0x5FE000, 469766144, 0x5FE000) )
    {
        printf("img cpy E\n");
    }
    else
    {
        v10.vrl_start = 0xBFC20000;
        v10.img_start = 0x1C000000;
        result = sub_3309C(&v10) + 1;
        if ( result )
            return result;
    }
}

```

load_fastboot function

Secureboot bypass: Fastboot stage

```
hwdog_certify->lock_state.fb_lock_stat = fb_lock_stat;
hwdog_certify->lock_state.user_lock_stat = 0;

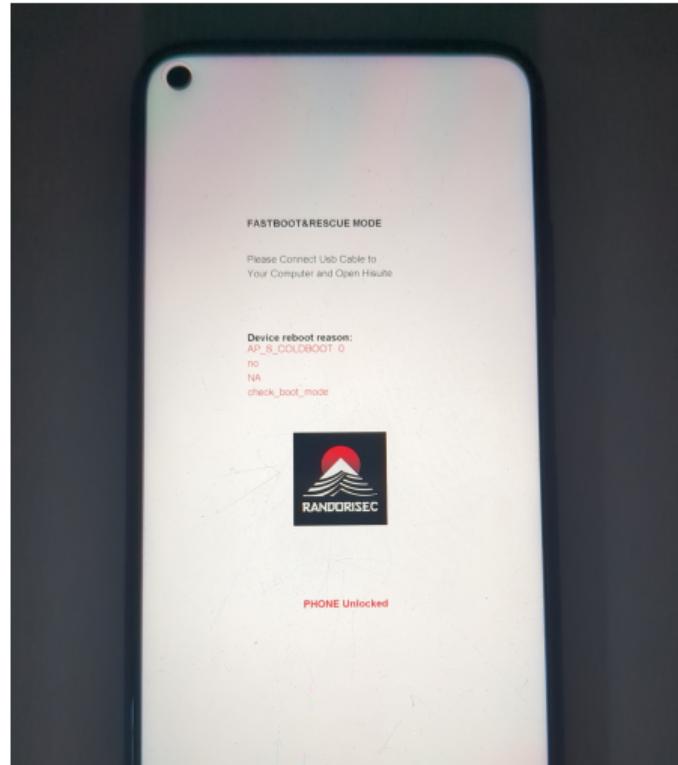
fb_lock_stat = get_hwdog_certify()->lock_state.fb_lock_stat;
if ( fb_lock_stat == 1 )
{
    fb_lock_stat_str = "UNLOCKED";
}
else if ( fb_lock_stat )
{
    fb_lock_stat_str = "UNKNOWN";
}
else
{
    fb_lock_stat_str = "LOCKED";
}
```

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Now let's have fun !!!!

Questions ?



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