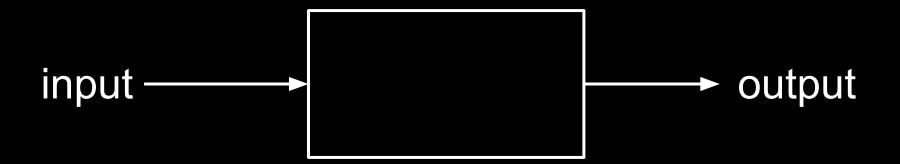
ALGORITHMS

what solves the problem?





algorithm, program, process

"A finitely long rule consisting of individual instructions is called an algorithm."

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf

"A program is an algorithm expressed in a programming language."

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf

"A process is a program that is currently executed by a computer."

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf



greatest common divisor

euclidean algorithm

Identify the larger number. If a < b, swap numbers so
that a > b

Repeat until one of the numbers becomes 0

Subtract b from a and replace a with the result

Return the number that is not zero

Loop 1: a = 18, $b = 48 \rightarrow swap$

Loop 1: a = 18, $b = 48 \rightarrow swap \rightarrow a = 48$, b = 18 a = 48 - 18 = 30

```
a = 18, b = 48 \rightarrow \text{swap} \rightarrow a = 48, b = 18

a = 48 - 18 = 30
```

Loop 2: a = 30, $b = 18 \rightarrow no swap$ a = 30 - 18 = 12

Loop 1:

```
Loop 1:
a = 18, b = 48 \rightarrow swap \rightarrow a = 48, b = 18
a = 48 - 18 = 30
Loop 2:
a = 30, b = 18 \rightarrow \text{no swap}
```

 $a = 12, b = 18 \rightarrow swap \rightarrow a = 18, b = 12$

a = 30 - 18 = 12

a = 18 - 12 = 6

Loop 3:

```
Loop 1:
a = 18, b = 48 \rightarrow swap \rightarrow a = 48, b = 18
a = 48 - 18 = 30
Loop 2:
a = 30, b = 18 \rightarrow no swap
```

Loop 3: a = 12, $b = 18 \rightarrow swap \rightarrow a = 18$, b = 12a = 18 - 12 = 6

a = 30 - 18 = 12

Loop 4:

$$a = 6, b = 12 \rightarrow \text{swap} \rightarrow a = 12, b = 6$$

 $a = 6, b = 12 \rightarrow swap \rightarrow a = 12, b = 6$ a = 12 - 6 = 6

```
Loop 1:
a = 18, b = 48 \rightarrow swap \rightarrow a = 48, b = 18
a = 48 - 18 = 30
Loop 2:
a = 30, b = 18 \rightarrow no swap
a = 30 - 18 = 12
```

Loop 3:

$$a = 12$$
, $b = 18 \rightarrow \text{swap} \rightarrow a = 18$, $b = 12$
 $a = 18 - 12 = 6$

Loop 4:

$$a = 6$$
, $b = 12 \rightarrow \text{swap} \rightarrow a = 12$, $b = 6$
 $a = 12 - 6 = 6$

Loop 5:

$$a = 6$$
, $b = 6 \rightarrow no swap$
 $a = 6 - 6 = 0$

```
Loop 1:

a = 18, b = 48 \rightarrow \text{swap} \rightarrow a = 48, b = 18

a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 \rightarrow \text{no swap}

a = 30 - 18 = 12

Loop 3:

a = 12, b = 18 \rightarrow \text{swap} \rightarrow a = 18, b = 12

a = 18 - 12 = 6
```

a = 6, $b = 12 \rightarrow swap \rightarrow a = 12$, b = 6

Loop 4:

Loop 5:

a = 12 - 6 = 6

a = 6 - 6 = 0

return b = 6

a = 6, $b = 6 \rightarrow no swap$

algorithm representation

100 Good Cookies I teaspoon baking soda 1 cup white sugar I cup quick oatmeal 1 cup brown sugar 1 cup coconut 1 cup margarine 1 cup Rice Krispies 1 cup cooking oil I cup chopped walnuts 1099 3/2 cups Hour 1 teaspoon vanilla I teaspoon cream of tartar 1-1202, package mini Chocolate chips Mix in order given. Drop by teaspoonful onto

cookie sheet. Bake at 350° for 8-10 minutes.

pseudocode

```
READ a, b
```

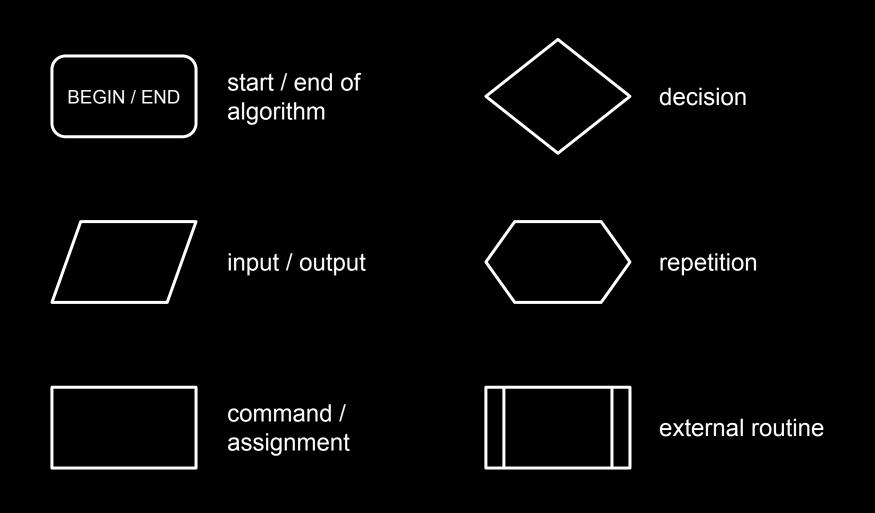
REPEAT

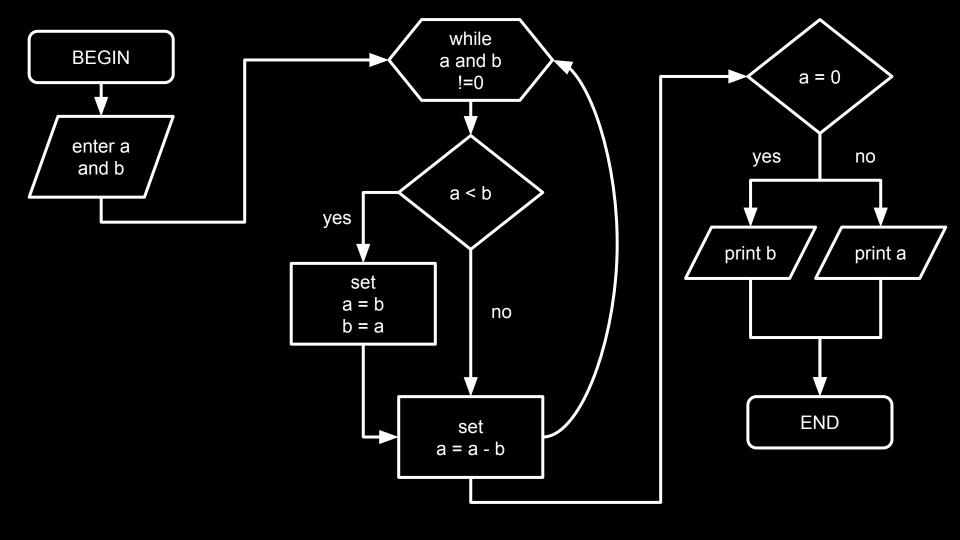
UNTIL
$$a = 0$$
 OR $b = 0$

a = a - b

RETURN the variable that is not 0

flow diagrams







square roots

babylonian method

calculate square root of x = 16

$$A = 1$$

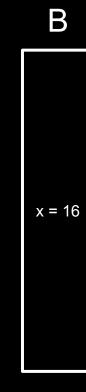
 $B = x / A = 16$

B

A x = 16

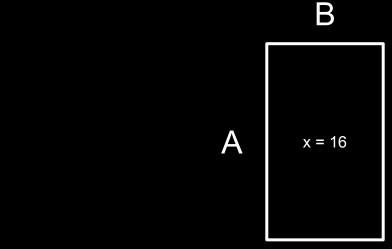
$$A = (A + B) / 2 = 17 / 2 = 8.5$$

 $B = x / A = 16 / 8.5 \approx 1.88$



$$A = (A + B) / 2 \approx 10.38 / 2 \approx 5.19$$

 $B = x / A \approx 16 / 5.19 \approx 3.08$

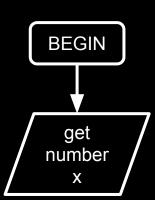


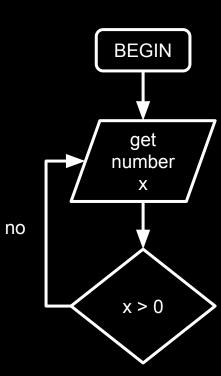
$$A = (A + B) / 2 \approx 8.27 / 2 \approx 4.14$$

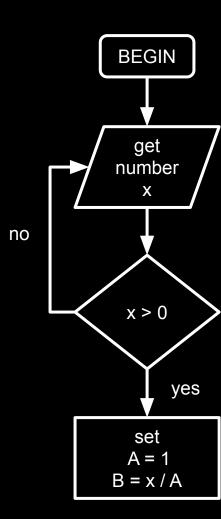
 $B = x / A \approx 16 / 4.14 \approx 3.86$

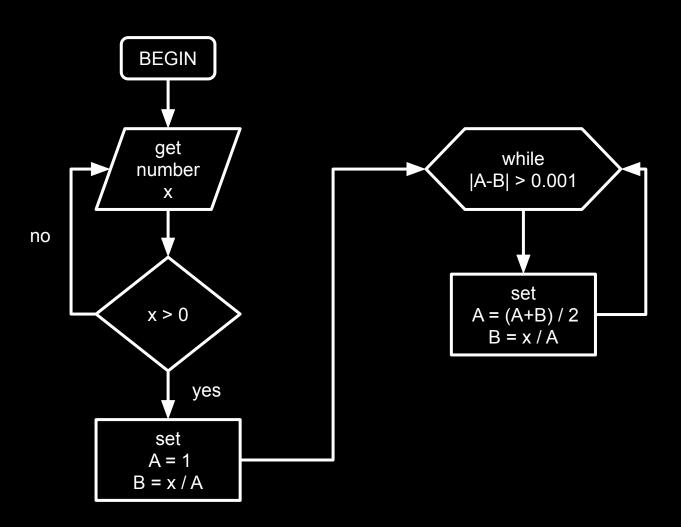
$$A = (A + B) / 2 = 8 / 2 = 4$$

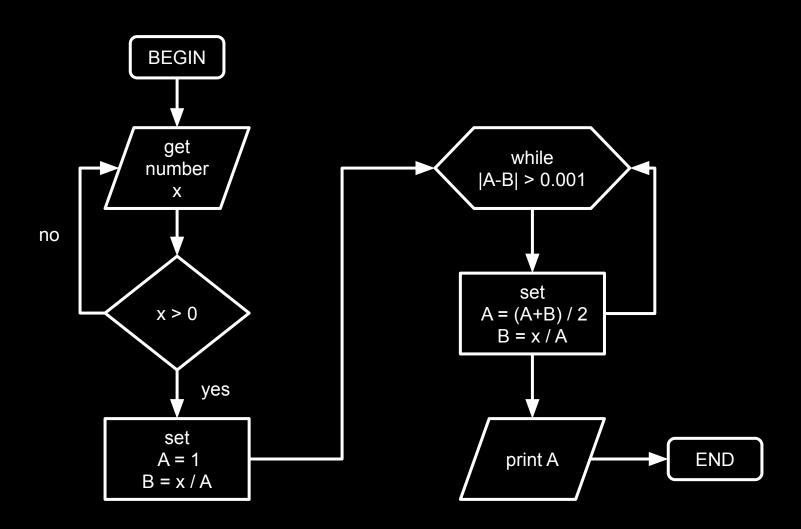
 $B = x / A = 16 / 4 = 4$





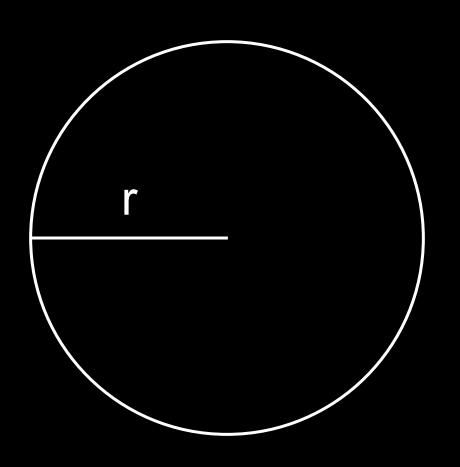


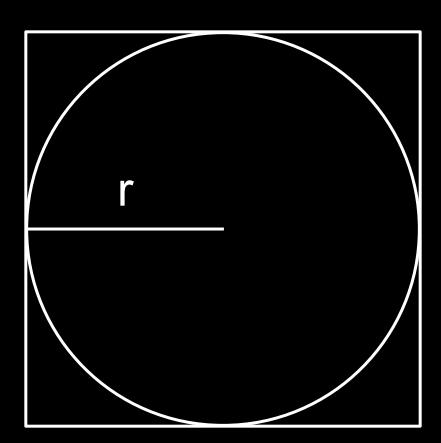


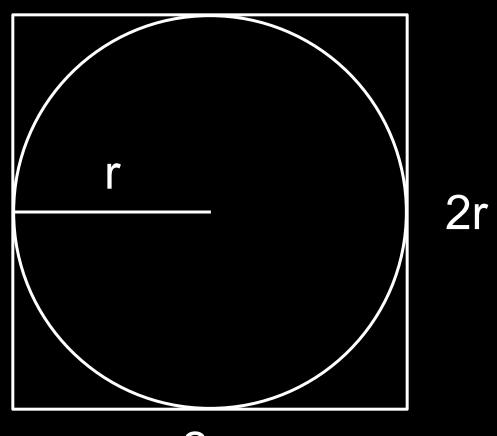




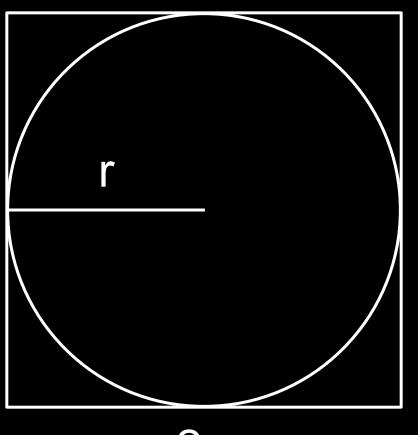
estimating π



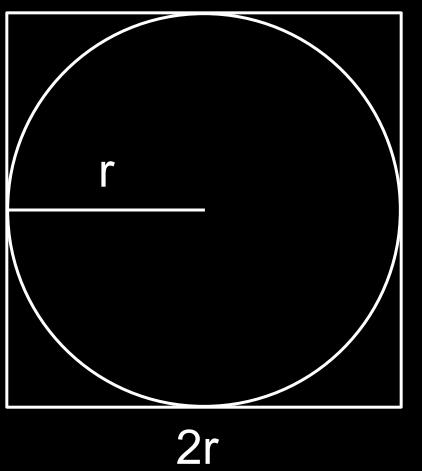




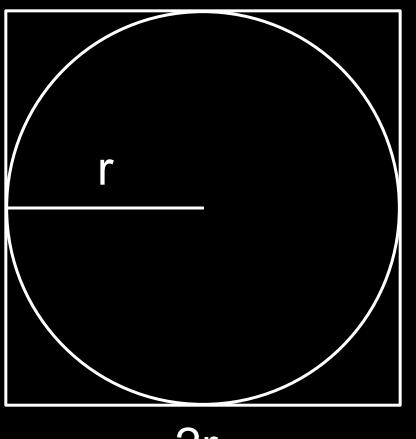
<u>2r</u>



$$2r \qquad \frac{\bigcirc}{\square} = \frac{\pi r^2}{4r^2}$$

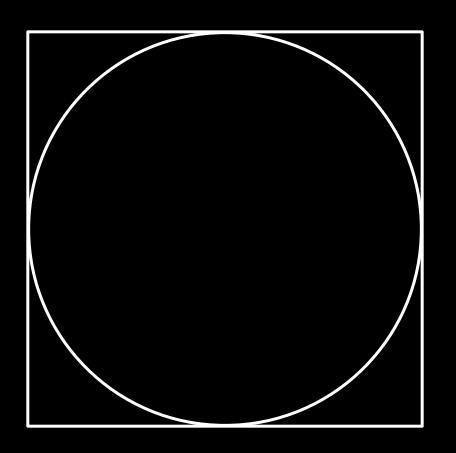


$$2r \qquad \frac{\bigcirc}{\square} = \frac{\pi r^2}{4r^2}$$

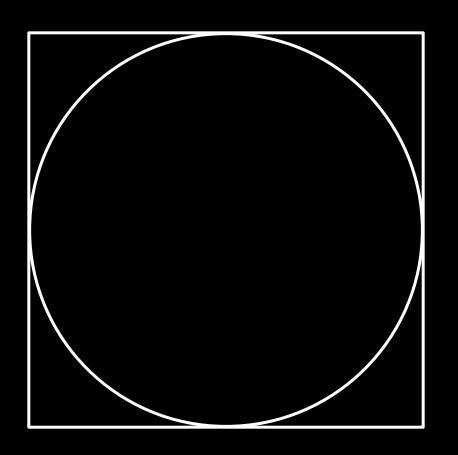


$$2r \qquad \frac{\bigcirc}{\square} = \frac{\pi}{4}$$

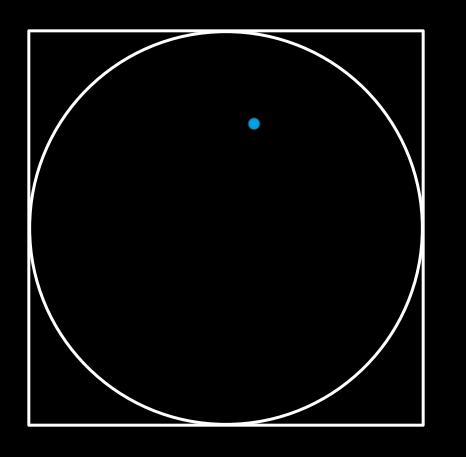
monte carlo simulation



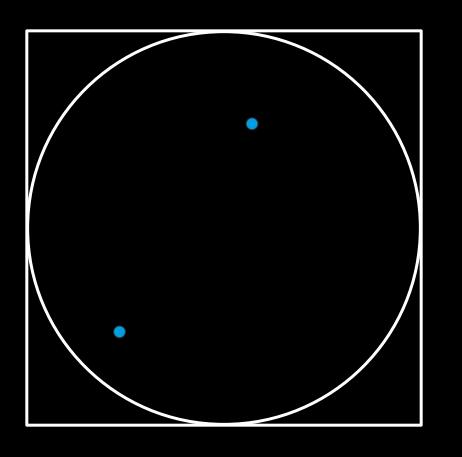
$$\frac{\bigcirc}{\square} = \frac{\pi}{4}$$



$$4 \quad \frac{\bigcirc}{\square} \quad = \quad \pi$$

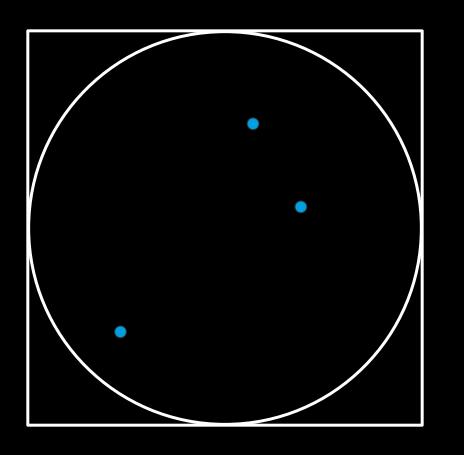


$$4 \frac{\bigcirc}{\square} = \pi$$



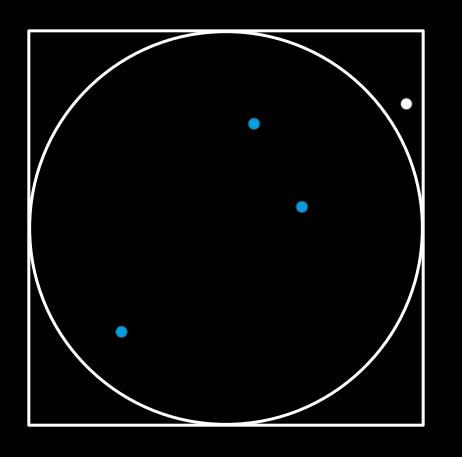
$$4 \frac{\bigcirc}{\square} = \pi$$

= 4

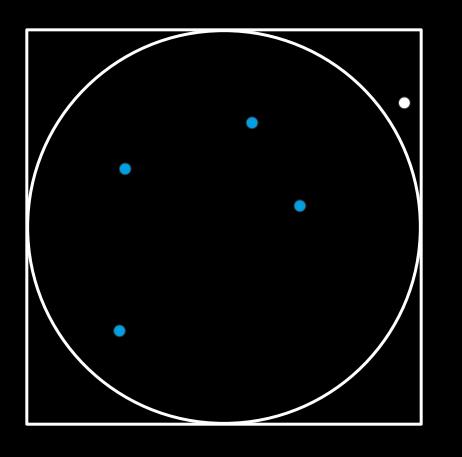


$$4 \frac{\bigcirc}{\square} = \pi$$

= 4

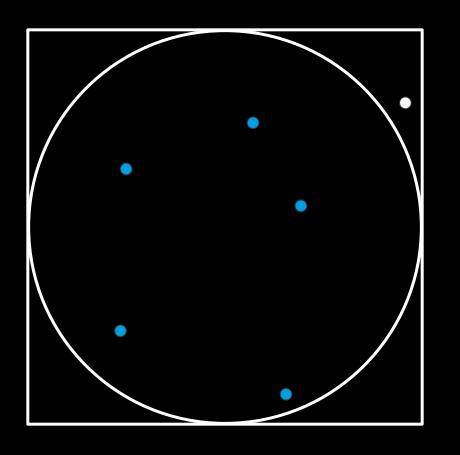


$$4 \frac{\bigcirc}{\square} = \pi$$

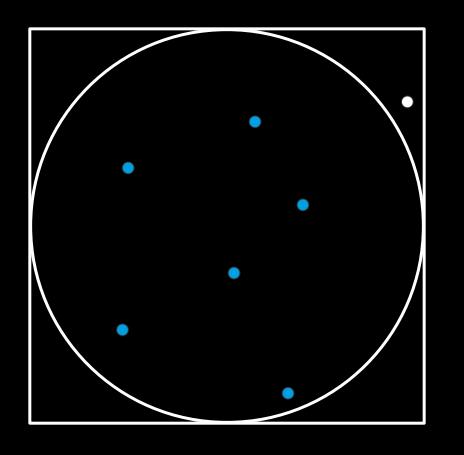


$$4 \frac{\bigcirc}{\square} = \pi$$

= 3,2

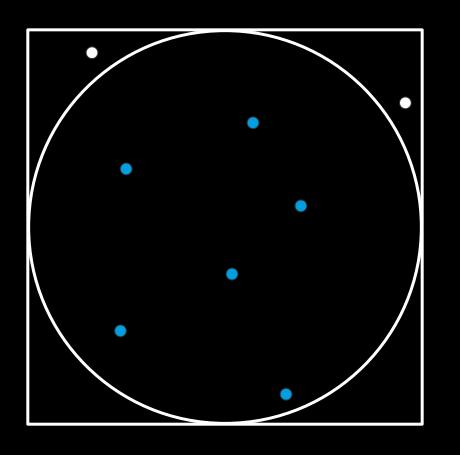


$$4 \frac{\bigcirc}{\square} = \pi$$

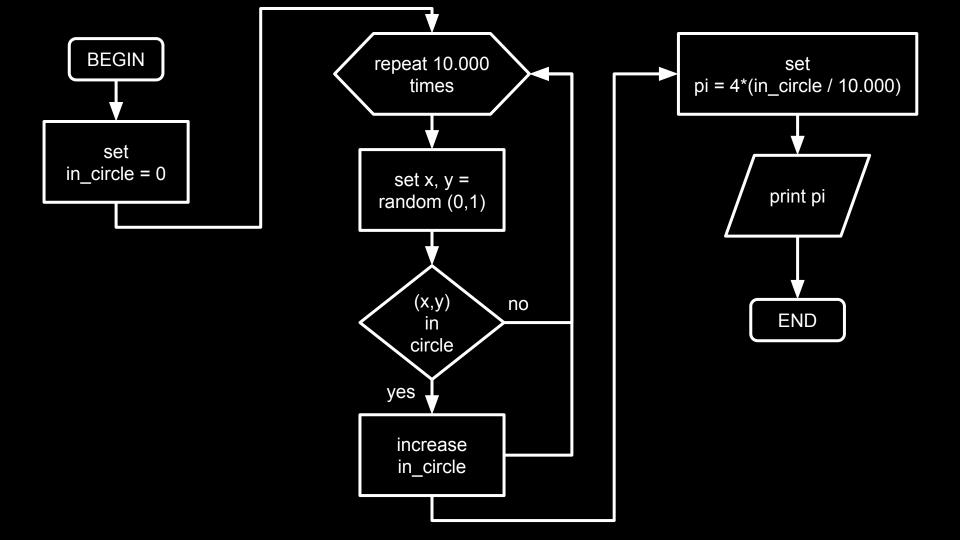


$$4 \frac{\bigcirc}{\square} = \pi$$

= 3,43



$$4 \frac{\bigcirc}{\square} = \pi$$



gregory-leibniz series

$$\pi = 4 \sum_{n=0}^{\infty} \frac{(-1)^n}{2n+1}$$

$$\pi = 4(1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \frac{1}{9} - \frac{1}{11} + \ldots)$$



sorting

[9, 5, 2, 1, 4, 7]

bubble sort

repeatedly compare and swap elements

until done.

[9, 5, 2, 1, 4, 7]

$$[9, 5, 2, 1, 4, 7] \longrightarrow 9 > 5? \xrightarrow{\text{yes}} [5, 9, 2, 1, 4, 7]$$

 $[5, 9, 2, 1, 4, 7] \longrightarrow 9 > 2? \xrightarrow{yes} [5, 2, 9, 1, 4, 7]$

$$[9, 5, 2, 1, 4, 7] \longrightarrow 9 > 5? \xrightarrow{\text{yes}} [5, 9, 2, 1, 4, 7]$$

$$[5, 9, 2, 1, 4, 7] \longrightarrow 9 > 2? \xrightarrow{\text{yes}} [5, 2, 9, 1, 4, 7]$$

 $[5, 2, 9, 1, 4, 7] \longrightarrow 9 > 1? \xrightarrow{\text{yes}} [5, 2, 1, 9, 4, 7]$

 $[5, 2, 1, 9, 4, 7] \longrightarrow 9 > 4? \xrightarrow{\text{yes}} [5, 2, 1, 4, 9, 7]$ $[5, 2, 1, 4, 9, 7] \longrightarrow 9 > 7? \xrightarrow{\text{yes}} [5, 2, 1, 4, 7, 9]$

$$[5, 2, 1, 4, 7, 9] \longrightarrow 5 > 2? \xrightarrow{\text{yes}} [2, 5, 1, 4, 7, 9]$$

$$[5, 2, 1, 4, 7, 9] \longrightarrow 5 > 2? \xrightarrow{\text{yes}} [2, 5, 1, 4, 7, 9]$$

 $[2, 1, 5, 4, 7, 9] \longrightarrow 5 > 4? \xrightarrow{yes} [2, 1, 4, 5, 7, 9]$

 $[2, 5, 1, 4, 7, 9] \longrightarrow 5 > 1? \xrightarrow{\text{yes}} [2, 1, 5, 4, 7, 9]$

$$[5, 2, 1, 4, 7, 9] \longrightarrow 5 > 2? \xrightarrow{\text{yes}} [2, 5, 1, 4, 7, 9]$$

 $[2, 5, 1, 4, 7, 9] \longrightarrow 5 > 1? \xrightarrow{\text{yes}} [2, 1, 5, 4, 7, 9]$

 $[2, 1, 5, 4, 7, 9] \longrightarrow 5 > 4? \xrightarrow{\text{yes}} [2, 1, 4, 5, 7, 9]$

 $[2, 1, 4, 5, 7, 9] \longrightarrow 5 > 7? \xrightarrow{\text{no}} [2, 1, 4, 5, 7, 9]$

 $[1, 2, 4, 5, 7, 9] \longrightarrow 2 > 4? \longrightarrow [1, 2, 4, 5, 7, 9]$

$$[1, 2, 4, 5, 7, 9] \longrightarrow 1 > 2? \xrightarrow{\text{no}} [1, 2, 4, 5, 7, 9]$$

 $[1, 2, 4, 5, 7, 9] \longrightarrow 2 > 4? \longrightarrow [1, 2, 4, 5, 7, 9]$

[1, 2, 4, 5, 7, 9] DONE!

selection sort

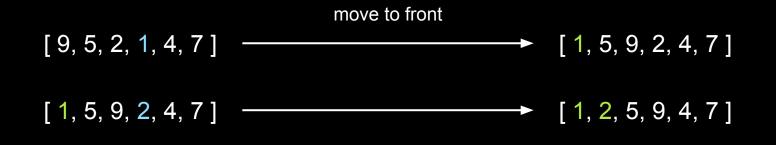
find the smallest element and move it to front. repeat for the rest of the elements.

move to front
[9, 5, 2, 1, 4, 7]

→ [1, 5, 9, 2, 4, 7]

move to front

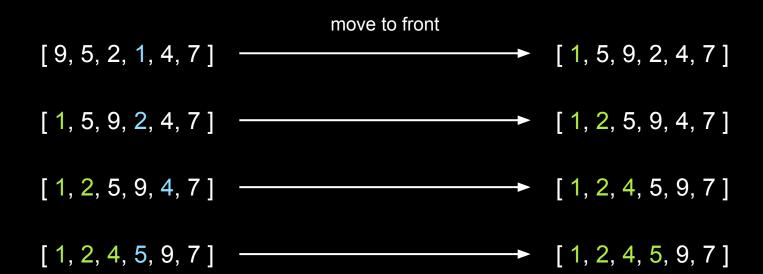
move to front [9, 5, 2, 1, 4, 7]

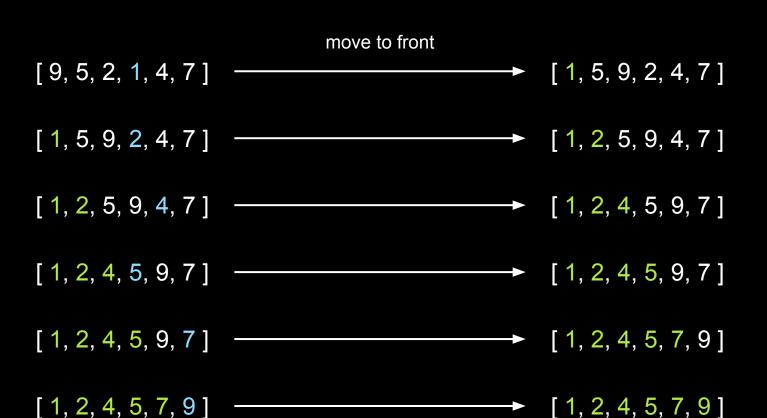


→ [1, 2, 4, 5, 9, 7]

[1, 2, 5, 9, 4, 7] ———

[1, 2, 4, 5, 9, 7]





merge sort

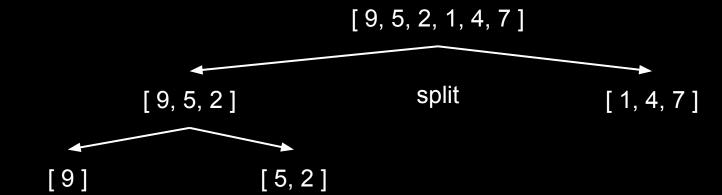
halves until only one element is left. then merge the sorted halves back together.

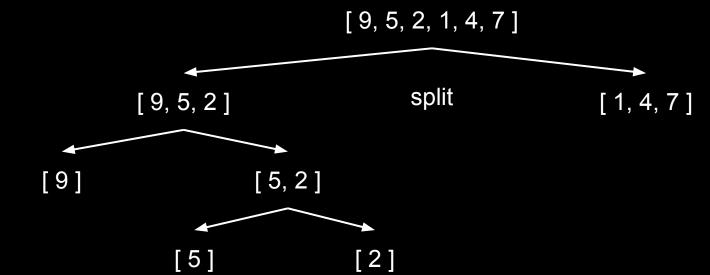
divide the elements recursively in two

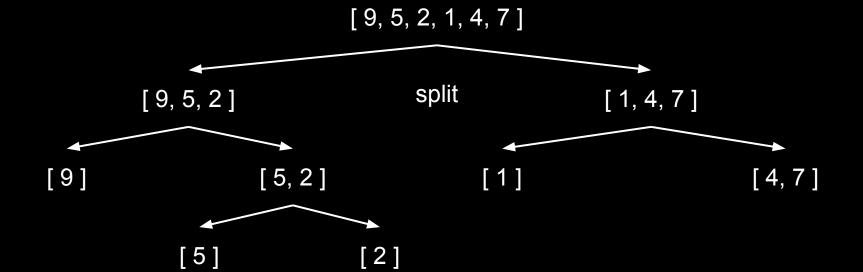
halves until only one element is left. then merge the sorted halves back together.

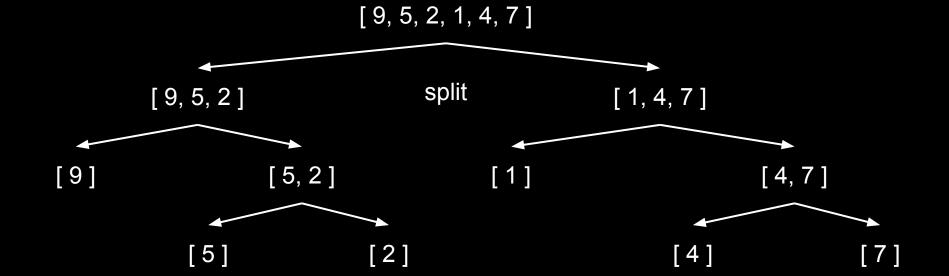
divide the elements recursively in two

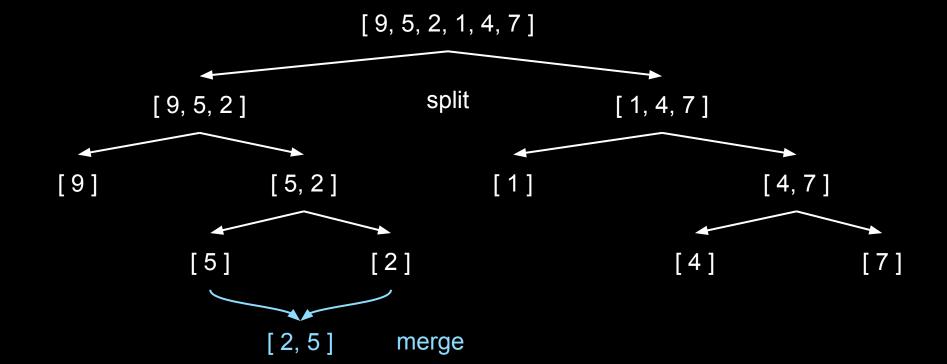


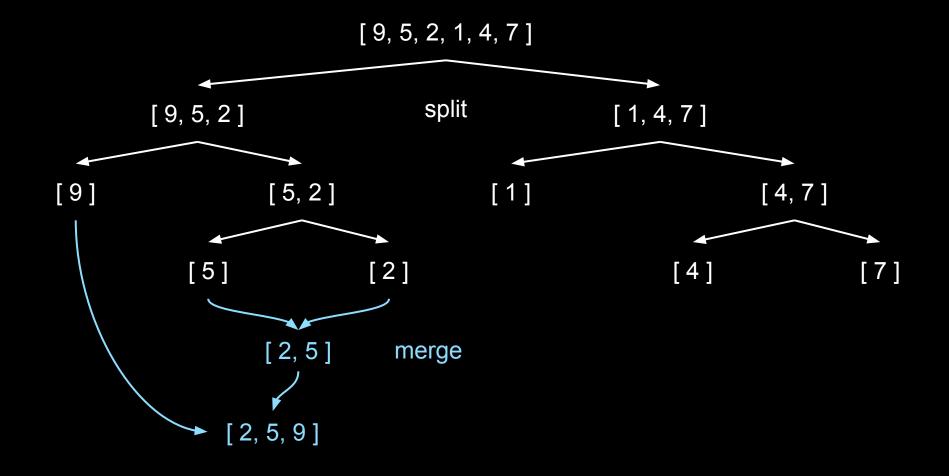


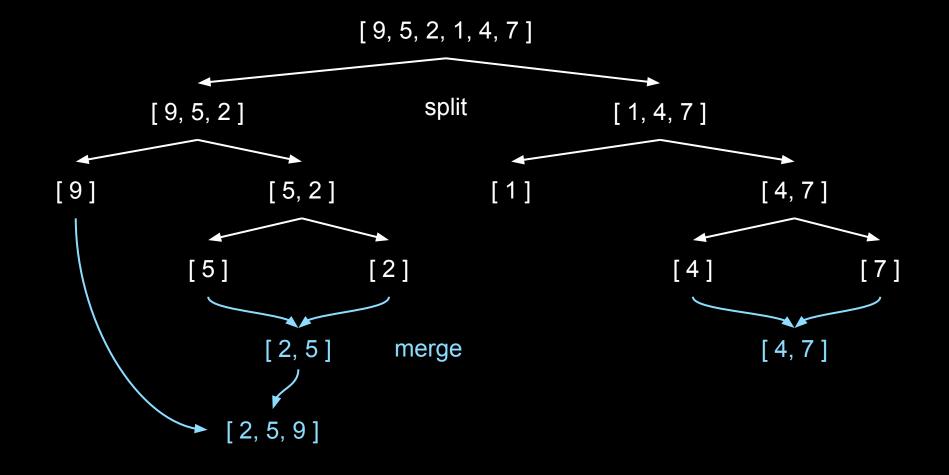


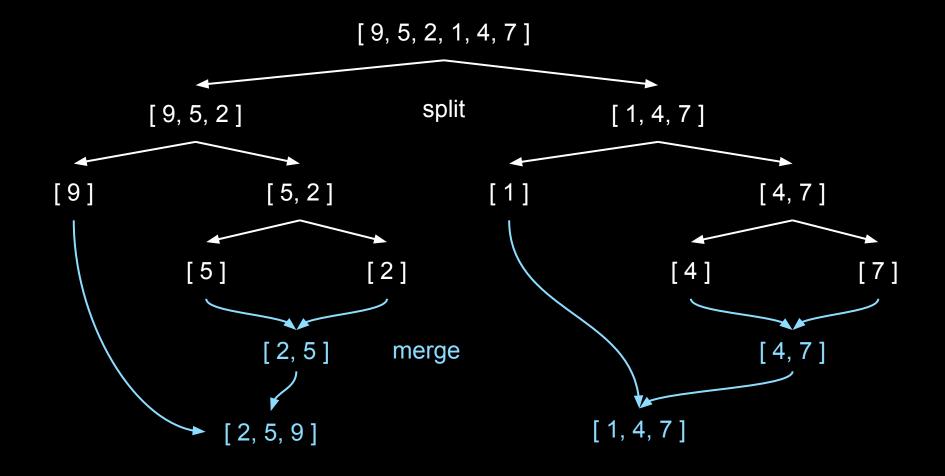


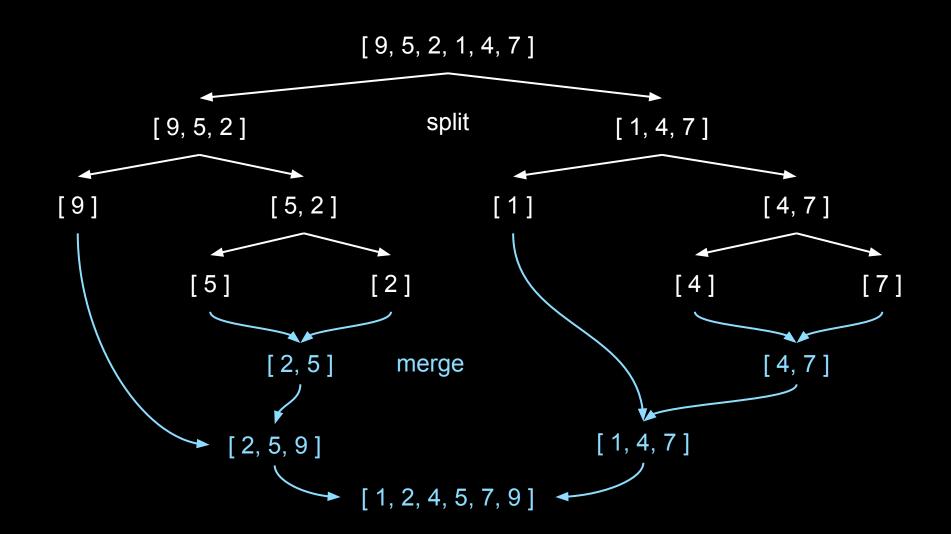




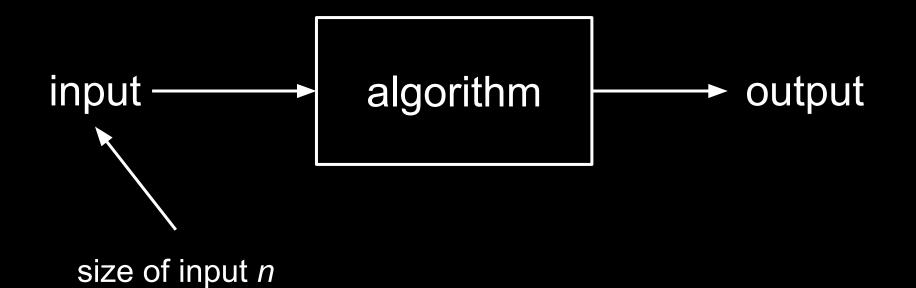


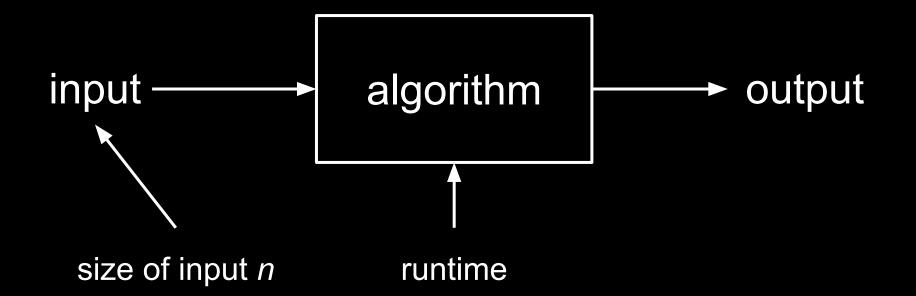




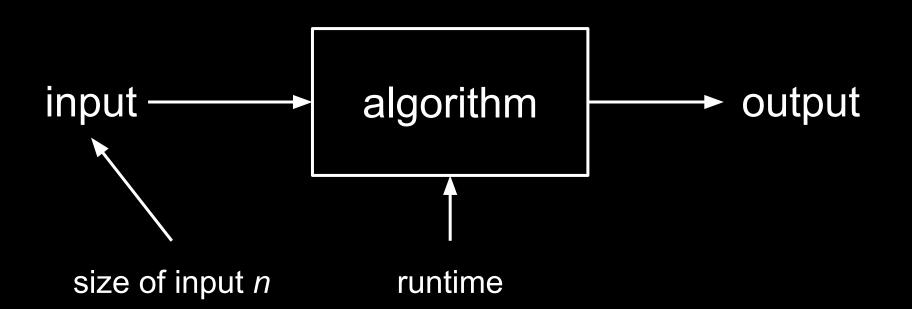


complexity





O(n)



runtime is constant and independent of problem size
runtime is determined by the logarithm of problem size
runtime is linear to problem size
runtime grows quadratically with the size of the problem
runtime grows cubically with the size of the problem
runtime grows exponentially with the size of the problem
runtime grows factorially with the size of the problem



optimization



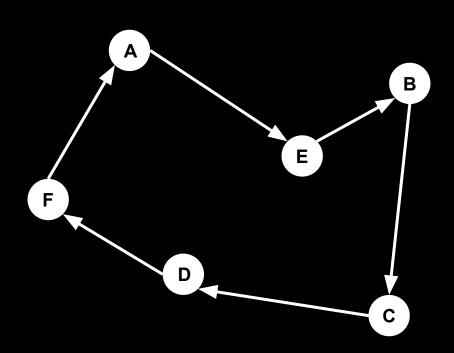
traveling salesmen

shortest tour?

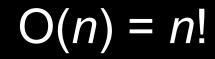
A B

D

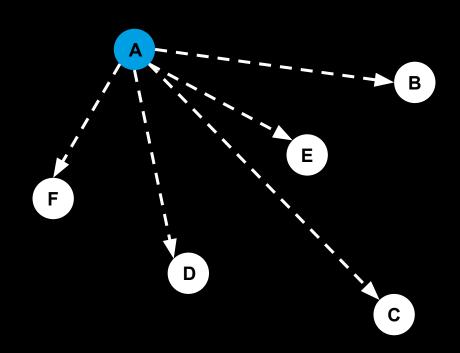
shortest tour?



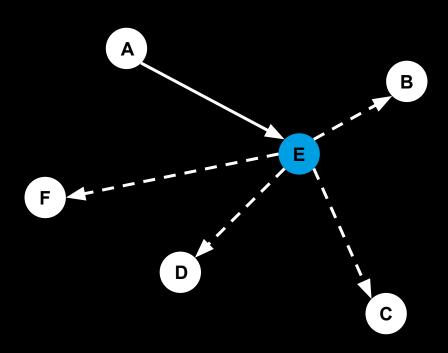
brute force



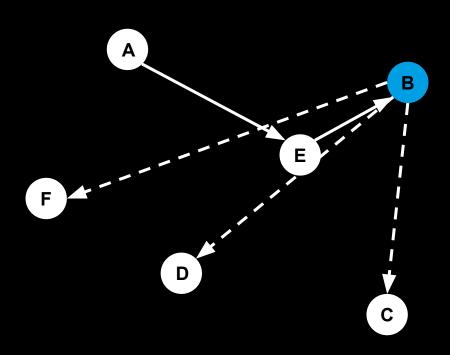
5 possible cities



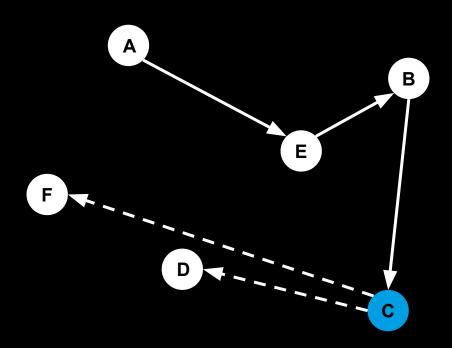
5 4 possible cities



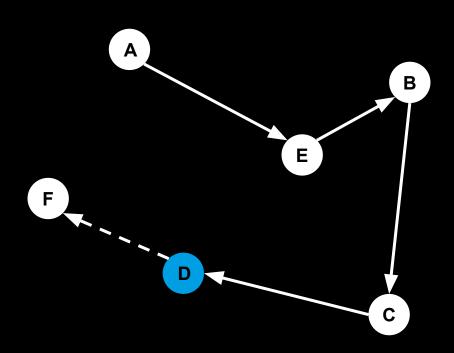
5 4 3 possible cities



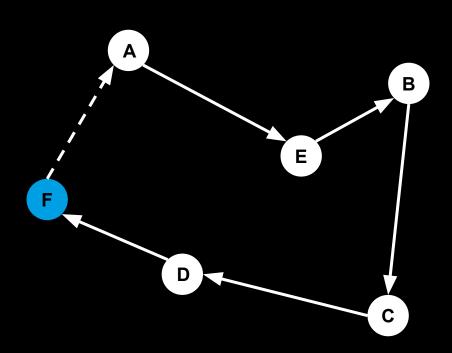
5 4 3 2 possible cities



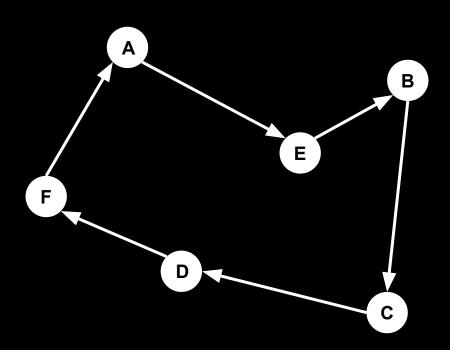
5 4 3 2 1 possible city



5 4 3 2 1 return nome



5 4 3 2 1 return nome



O(n) = n! n = 5

$$O(n) = n!$$
 $n = 5$

N = 5 * 4 * 3 * 2 * 1

$$O(n) = n!$$

$$n = 5$$

$$N = 5 * 4 * 3 * 2 * 1$$

$$= 120$$

n = 10 n = 20 n = 30

n = 25

brute force takes longer than the universe is old

n = 60

more possible routes than atoms in the universe

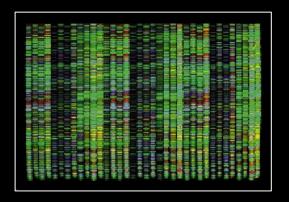






Image source: IEEE

Image source: VDI Nachrichten

Image source: Wikimedia

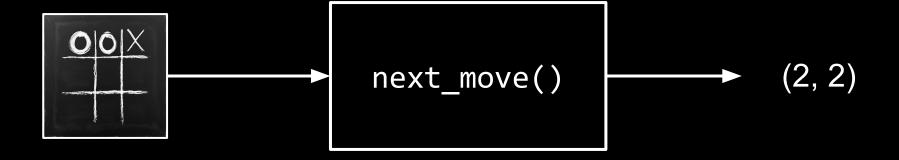






Image source: <u>IAS Observatory</u>

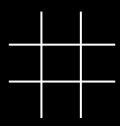
but sometimes, brute force works perfectly

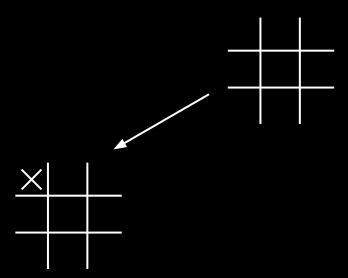


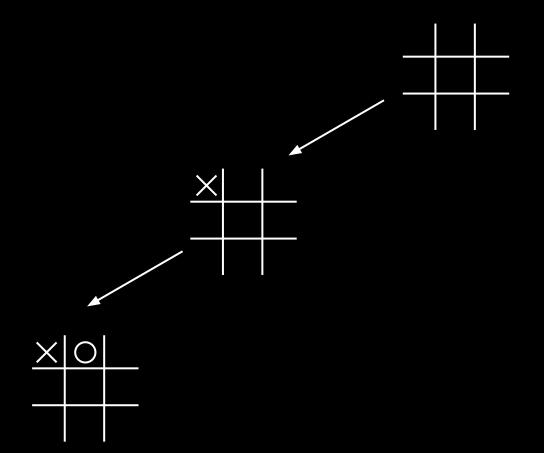
brute-force search:

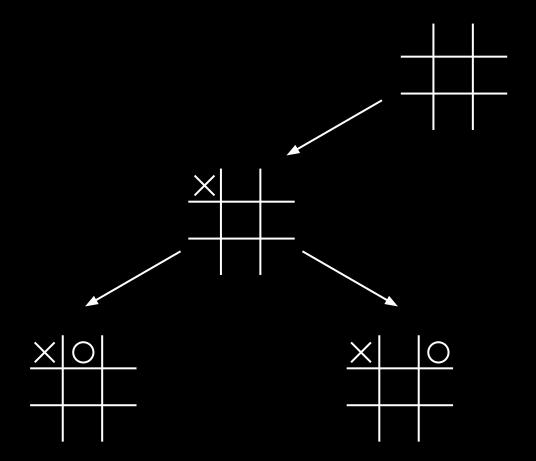
evaluate all possible move sequences.

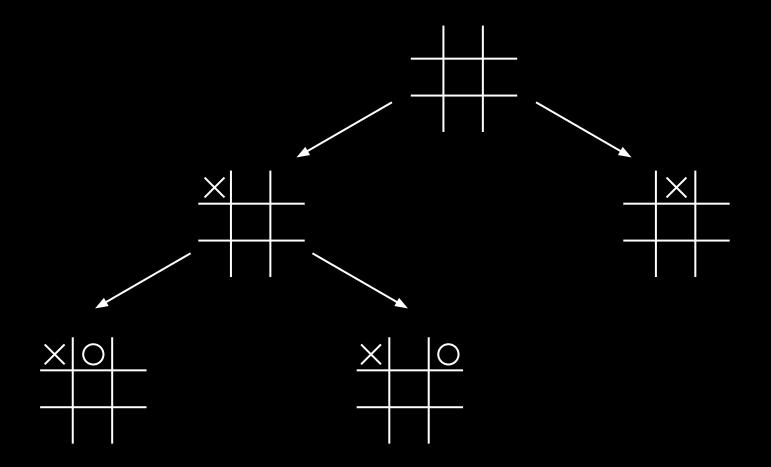
choose the move with the best value.

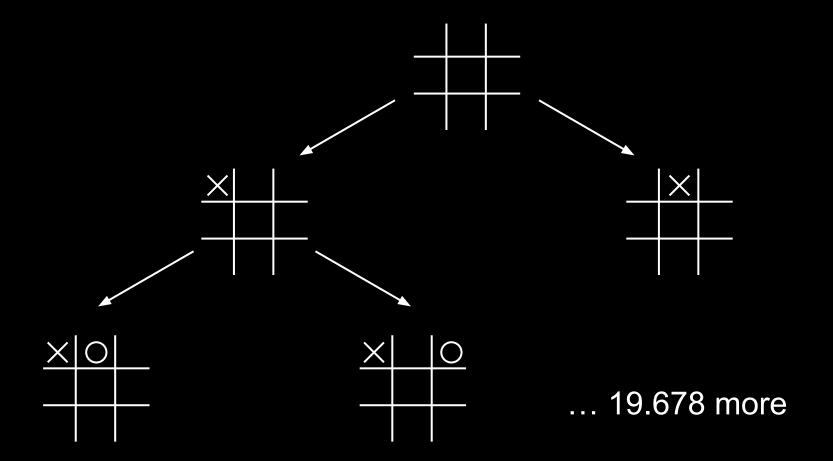






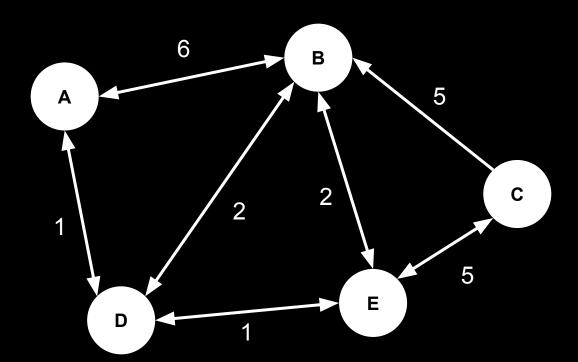






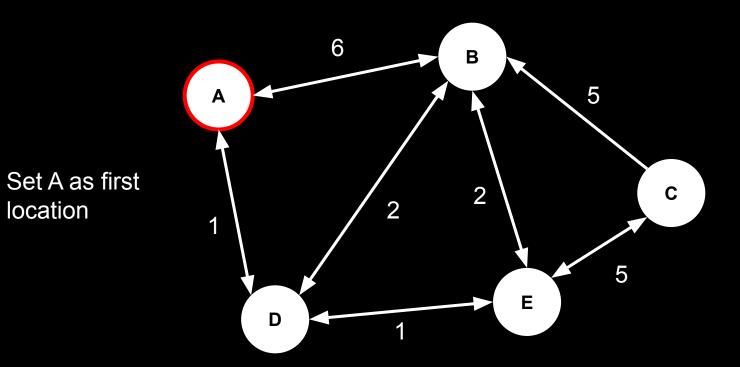


shortest paths

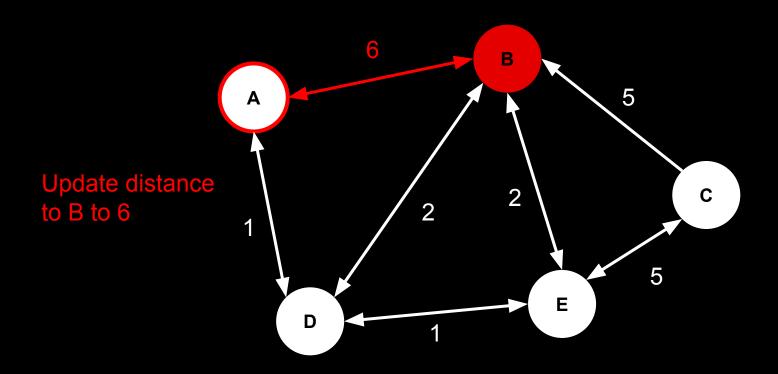


dijkstra's algorithm

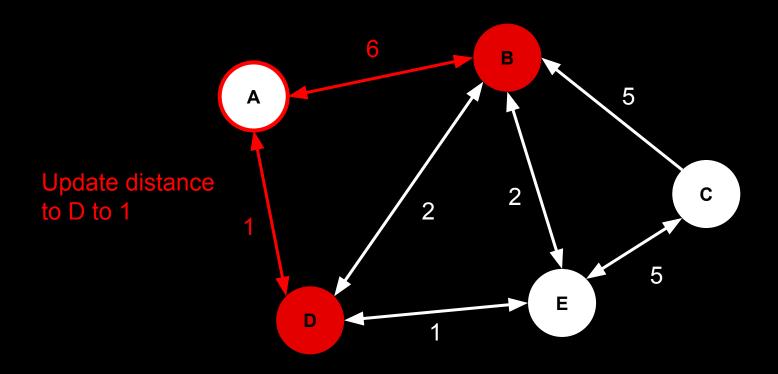
Distances: A = 0, $B = \infty$, $C = \infty$, $D = \infty$, $E = \infty$



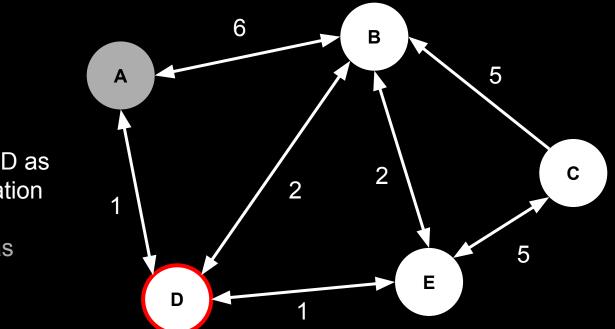
Distances: A = 0, B = 6, $C = \infty$, $D = \infty$, $E = \infty$



Distances: A = 0, B = 6, C = ∞, D = 1, E = ∞



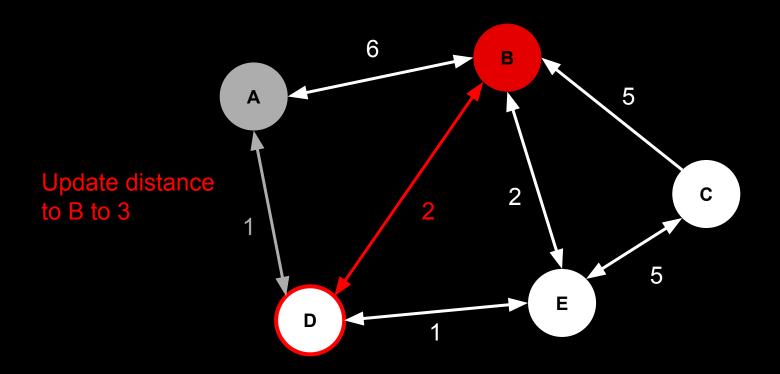
Distances: A = 0, B = 6, C = ∞, D = 1, E = ∞

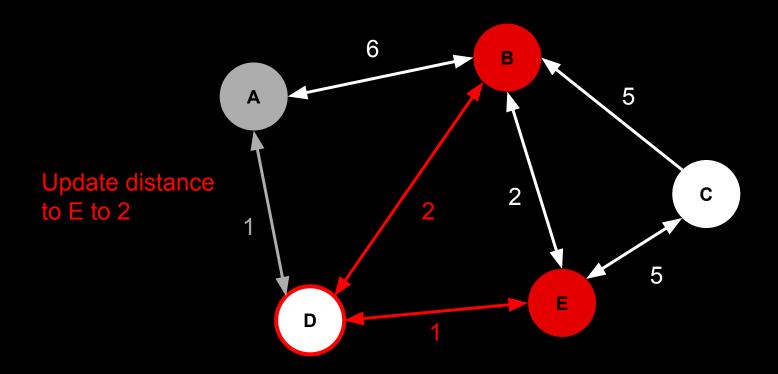


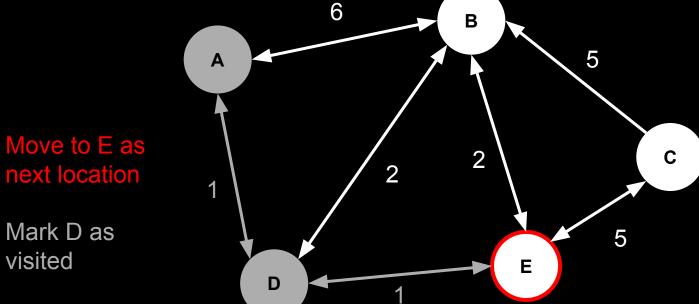
Move to D as next location

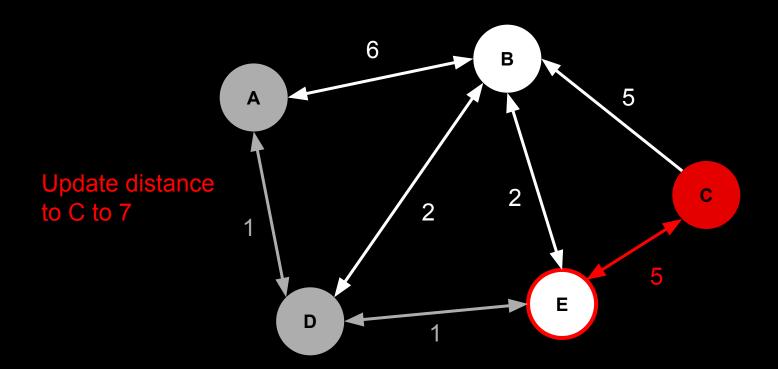
Mark A as visited

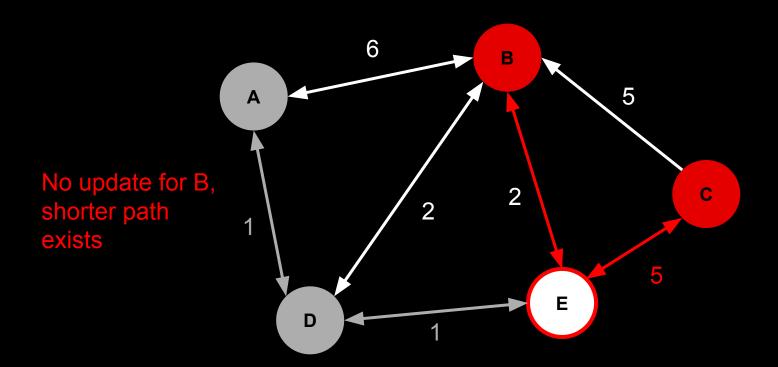
Distances: A = 0, B = 3, C = ∞, D = 1, E = ∞

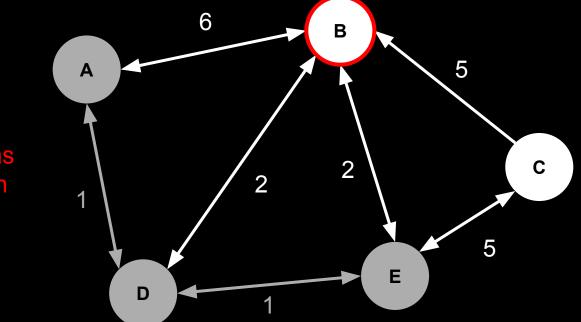












Move to B as new location

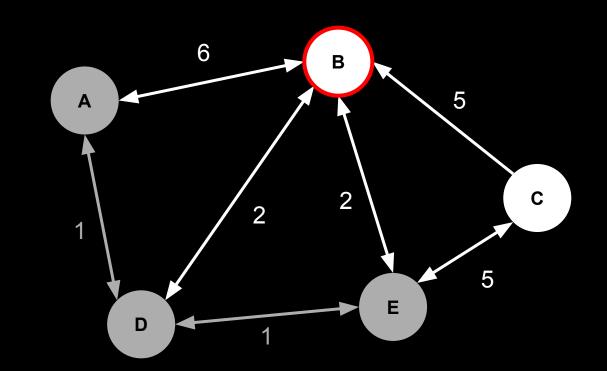
Mark E as visited

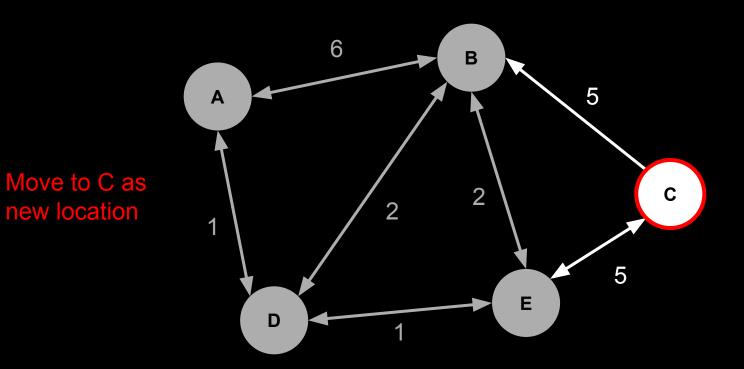
No further

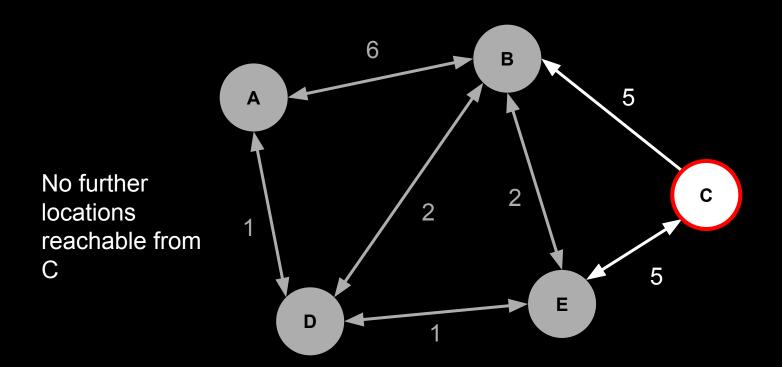
locations

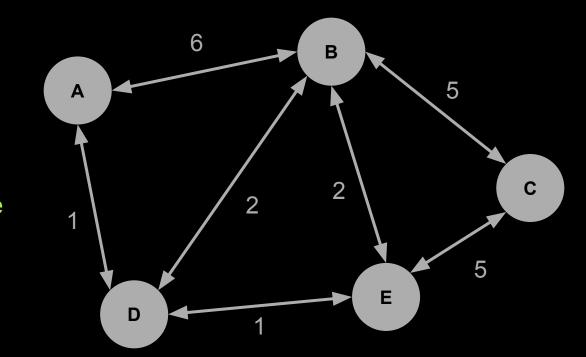
reachable

from B









All nodes visited, we're done!



spam emails



finding oranges in images

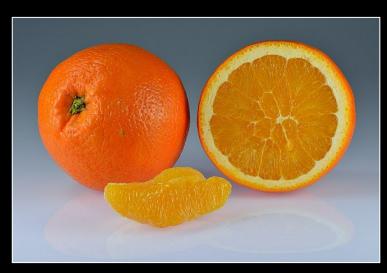


Image source: Wikimedia

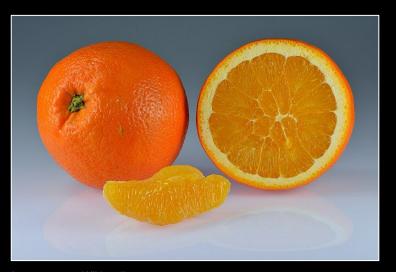


Image source: Wikimedia

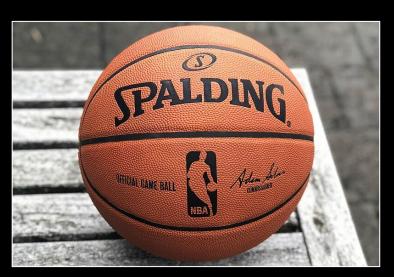


Image source: Wikimedia

what set of rules can solve this?

machine learning algorithms

