

# ALGORITHMS

[BACK](#)

# what solves the problem?





algorithm, program, process

"A finitely long rule consisting of individual instructions is called an **algorithm**."

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: <http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf>

"A **program** is an algorithm expressed in a  
programming language."

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: <http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf>

**"A process is a program that is currently executed by a computer."**

Source: Vornberger, O., Algorithmen und Datenstrukturen, Lecture notes: <http://www-lehre.inf.uos.de/~ainf/2013/PDF/skript.pdf>



# greatest common divisor

euclidean algorithm

Identify the larger number. If  $a < b$ , swap numbers so that  $a > b$

Subtract  $b$  from  $a$  and replace  $a$  with the result

Repeat until one of the numbers becomes 0

Return the number that is not zero

Loop 1:

a = 18, b = 48 → swap

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18

a = 48 - 18 = 30

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18  
a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 → no swap  
a = 30 - 18 = 12

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18  
a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 → no swap  
a = 30 - 18 = 12

Loop 3:

a = 12, b = 18 → swap → a = 18, b = 12  
a = 18 - 12 = 6

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18  
a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 → no swap  
a = 30 - 18 = 12

Loop 3:

a = 12, b = 18 → swap → a = 18, b = 12  
a = 18 - 12 = 6

Loop 4:

a = 6, b = 12 → swap → a = 12, b = 6  
a = 12 - 6 = 6

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18  
a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 → no swap  
a = 30 - 18 = 12

Loop 3:

a = 12, b = 18 → swap → a = 18, b = 12  
a = 18 - 12 = 6

Loop 4:

a = 6, b = 12 → swap → a = 12, b = 6  
a = 12 - 6 = 6

Loop 5:

a = 6, b = 6 → no swap  
a = 6 - 6 = 0

Loop 1:

a = 18, b = 48 → swap → a = 48, b = 18  
a = 48 - 18 = 30

Loop 2:

a = 30, b = 18 → no swap  
a = 30 - 18 = 12

Loop 3:

a = 12, b = 18 → swap → a = 18, b = 12  
a = 18 - 12 = 6

Loop 4:

a = 6, b = 12 → swap → a = 12, b = 6  
a = 12 - 6 = 6

Loop 5:

a = 6, b = 6 → no swap  
a = 6 - 6 = 0

return b = 6

# algorithm representation

## 100 Good Cookies

1 cup white sugar

1 cup brown sugar

1 cup margarine

1 cup cooking oil

1 egg

1 teaspoon vanilla

1 teaspoon cream of tartar

1 teaspoon baking soda

1 cup quick oatmeal

1 cup coconut

1 cup Rice Krispies

1 cup chopped walnuts

3½ cups flour

1-12oz. package mini

chocolate chips

Mix in order given. Drop by teaspoonful onto  
cookie sheet. Bake at 350° for 8-10 minutes.  
(over)

pseudocode

READ a, b

REPEAT

IF a < b THEN

a = b

b = a

a = a - b

UNTIL a = 0 OR b = 0

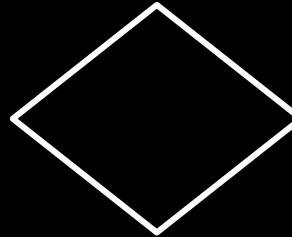
RETURN the variable that is not 0

# flow diagrams

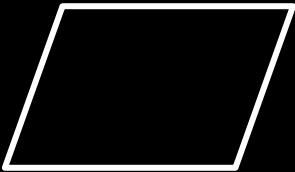


BEGIN / END

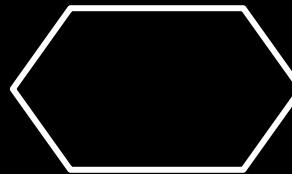
start / end of  
algorithm



decision



input / output



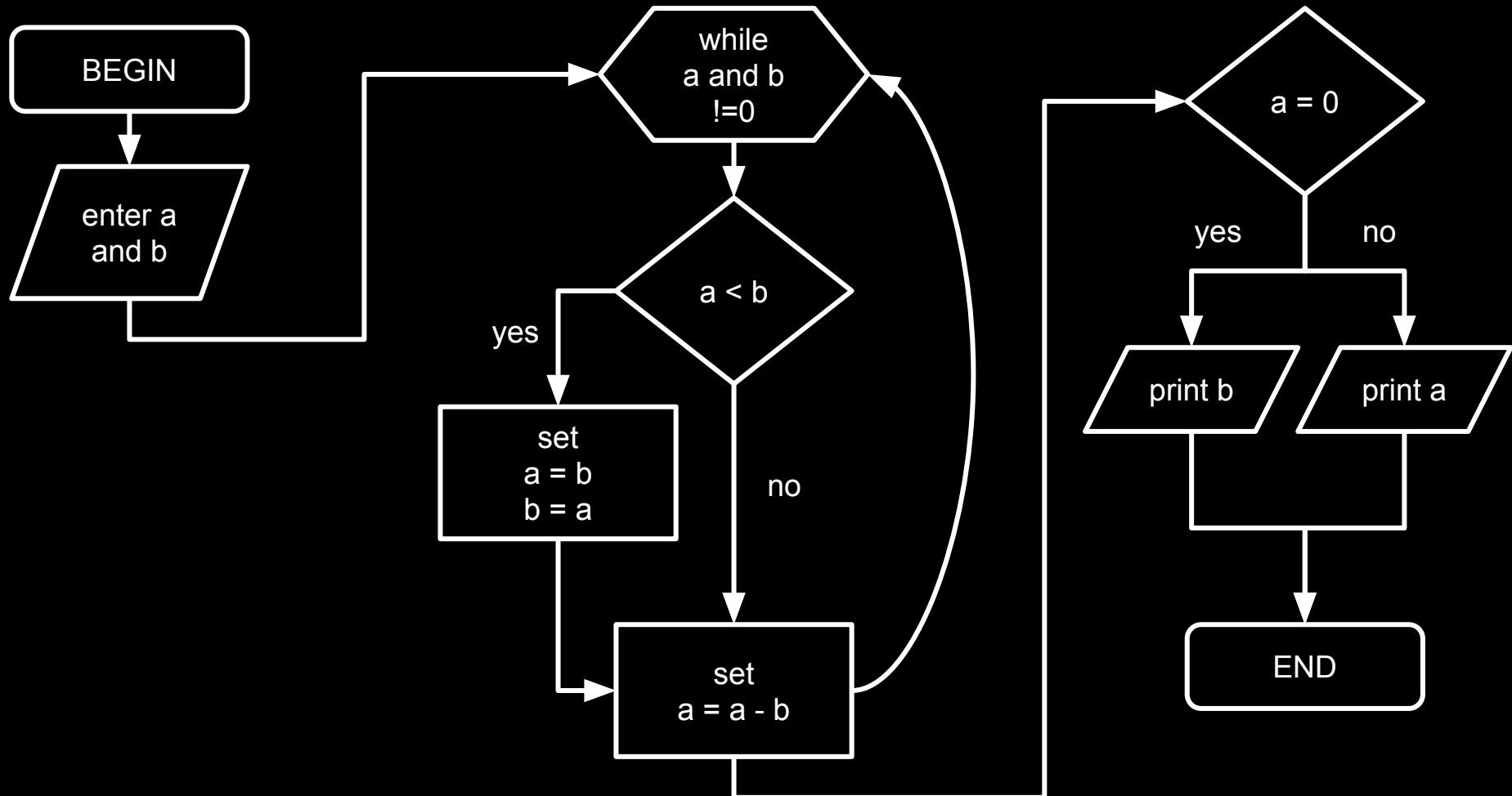
repetition



command /  
assignment



external routine





# square roots

# babylonian method

calculate square root of  
 $x = 16$

A = 1

B = x / A = 16

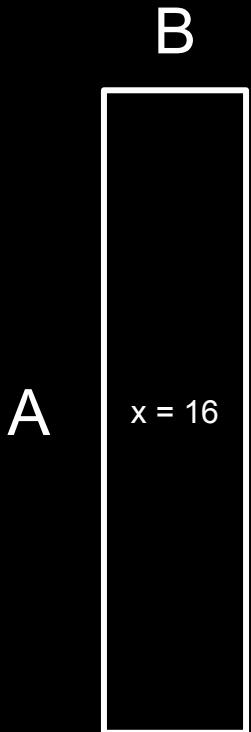
B

A

x = 16

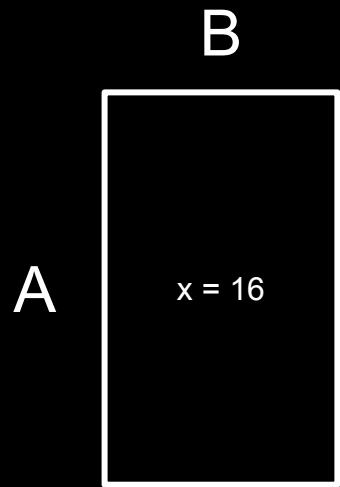
$$A = (A + B) / 2 = 17 / 2 = 8.5$$

$$B = x / A = 16 / 8.5 \approx 1.88$$



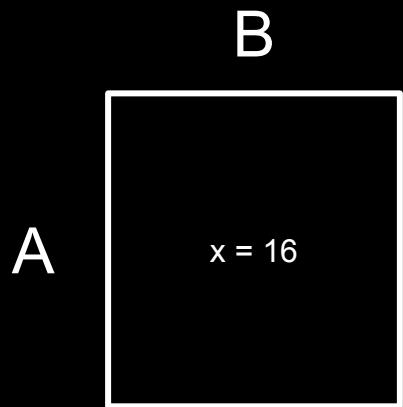
$$A = (A + B) / 2 \approx 10.38 / 2 \approx 5.19$$

$$B = x / A \approx 16 / 5.19 \approx 3.08$$



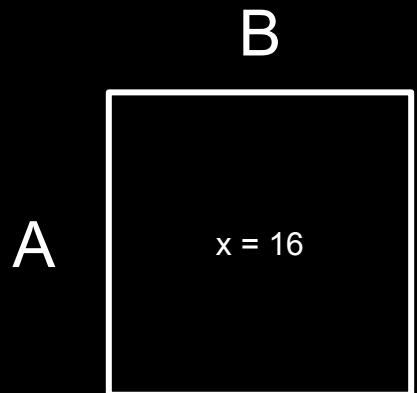
$$A = (A + B) / 2 \approx 8.27 / 2 \approx 4.14$$

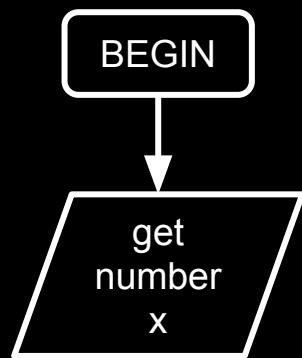
$$B = x / A \approx 16 / 4.14 \approx 3.86$$

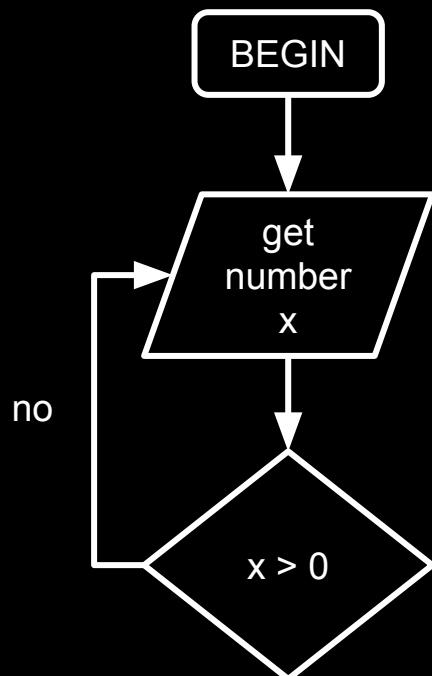


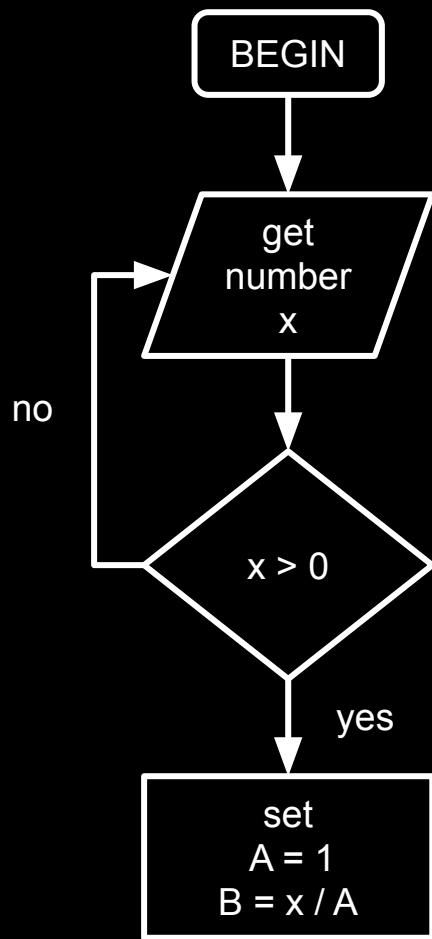
$$A = (A + B) / 2 = 8 / 2 = 4$$

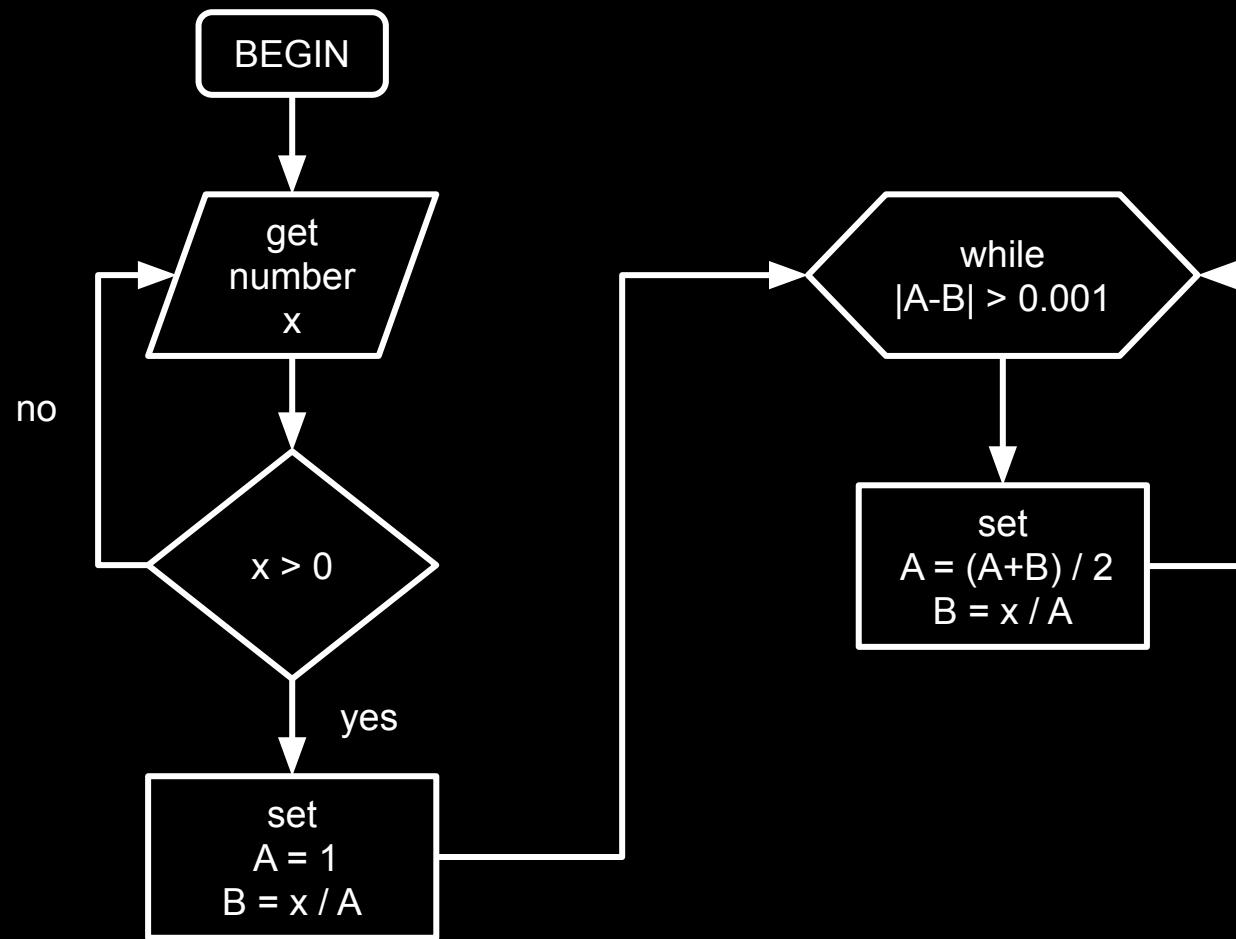
$$B = x / A = 16 / 4 = 4$$

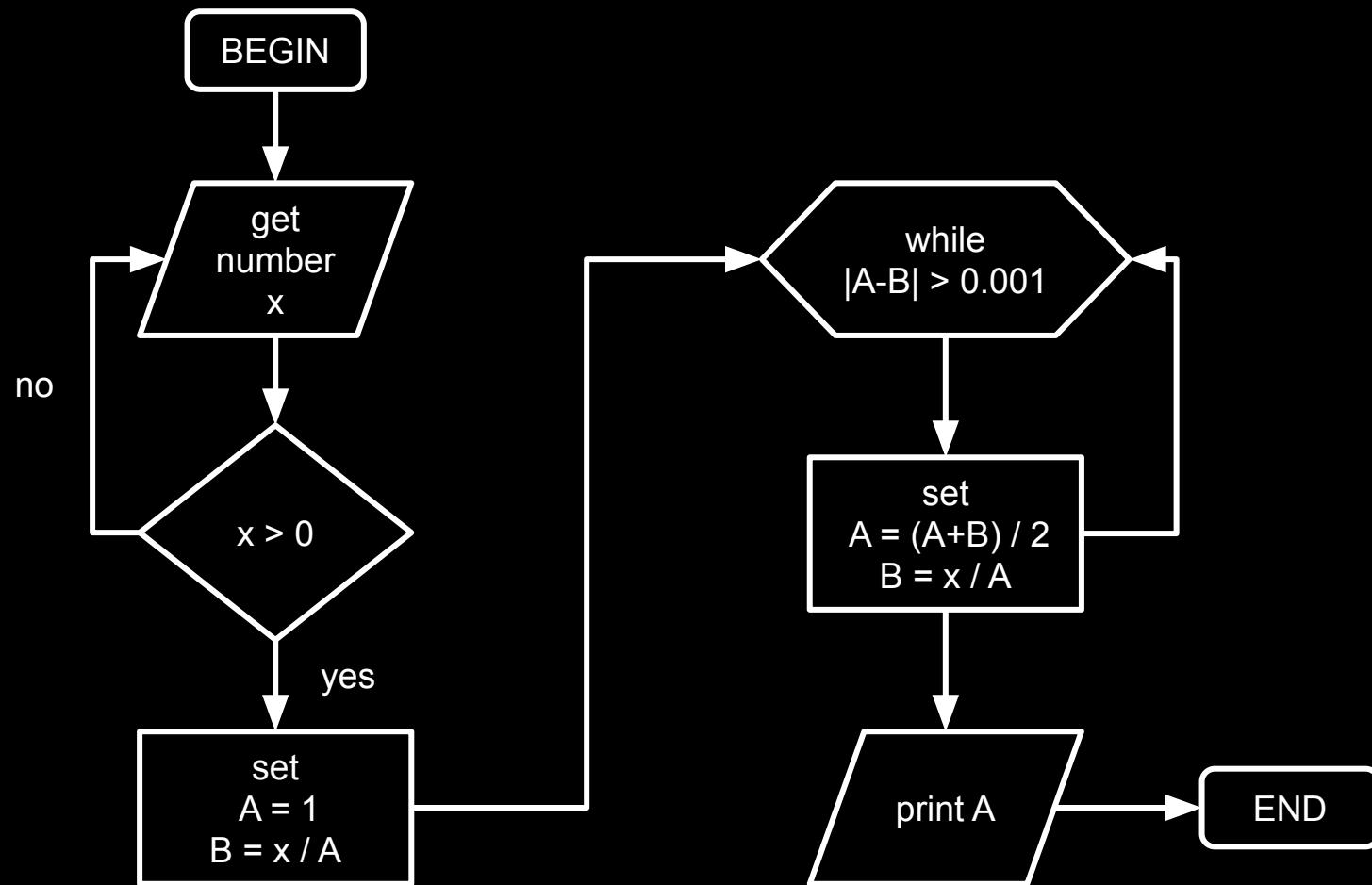






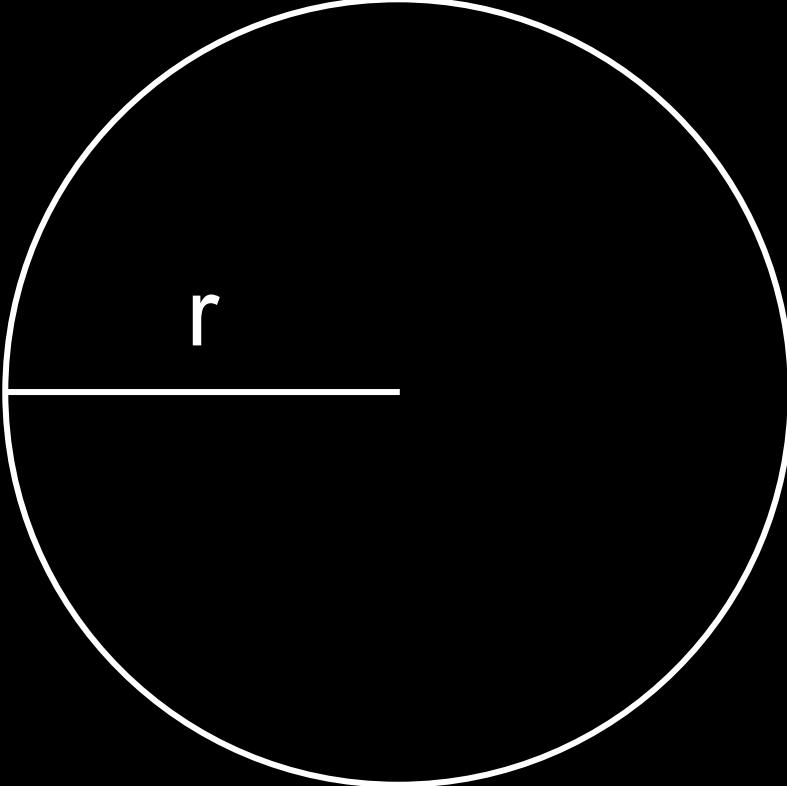






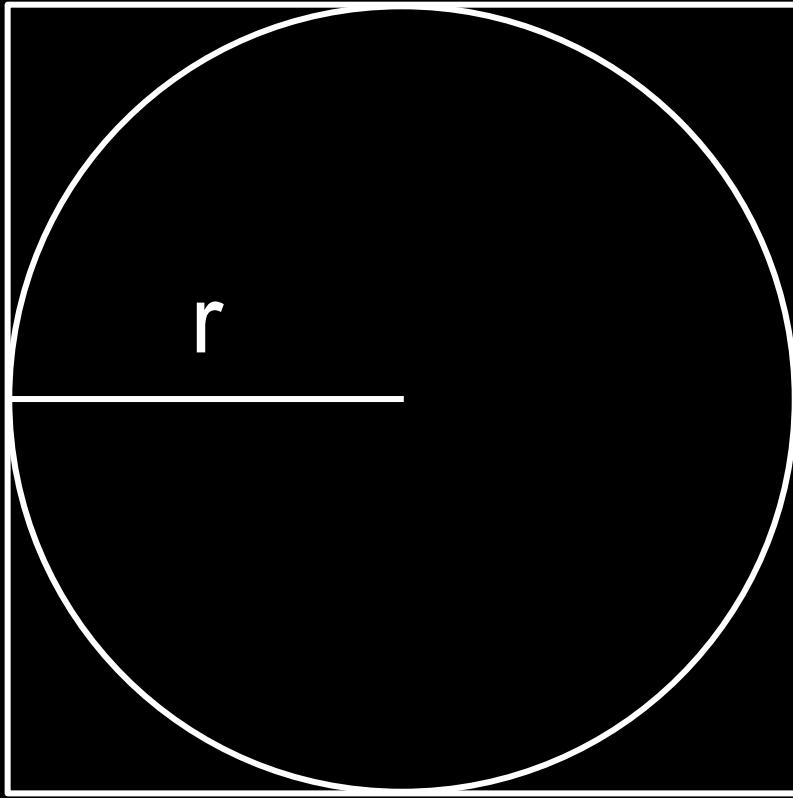


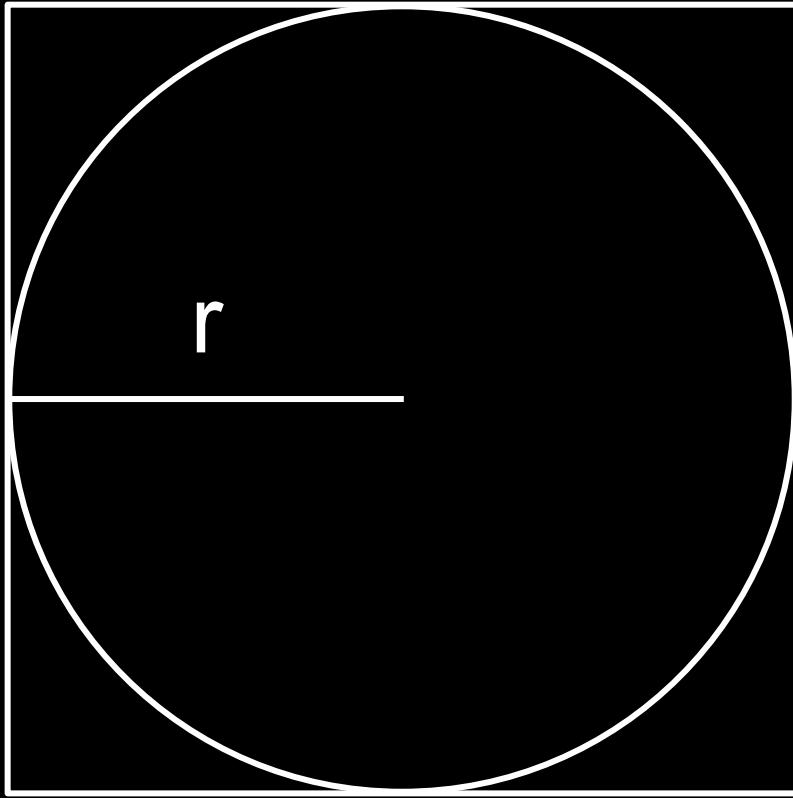
# estimating $\pi$

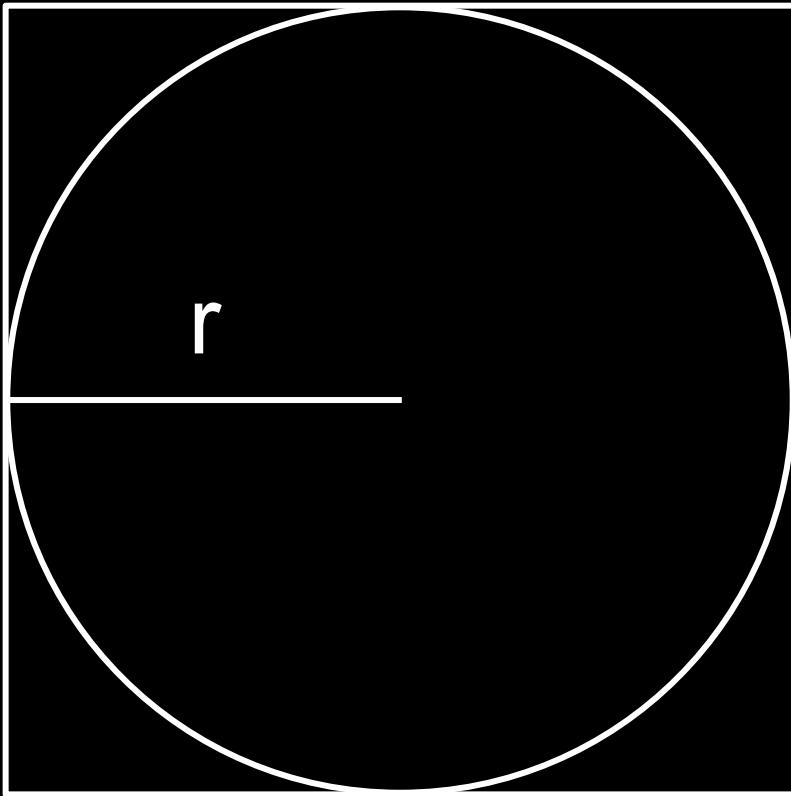


$r$

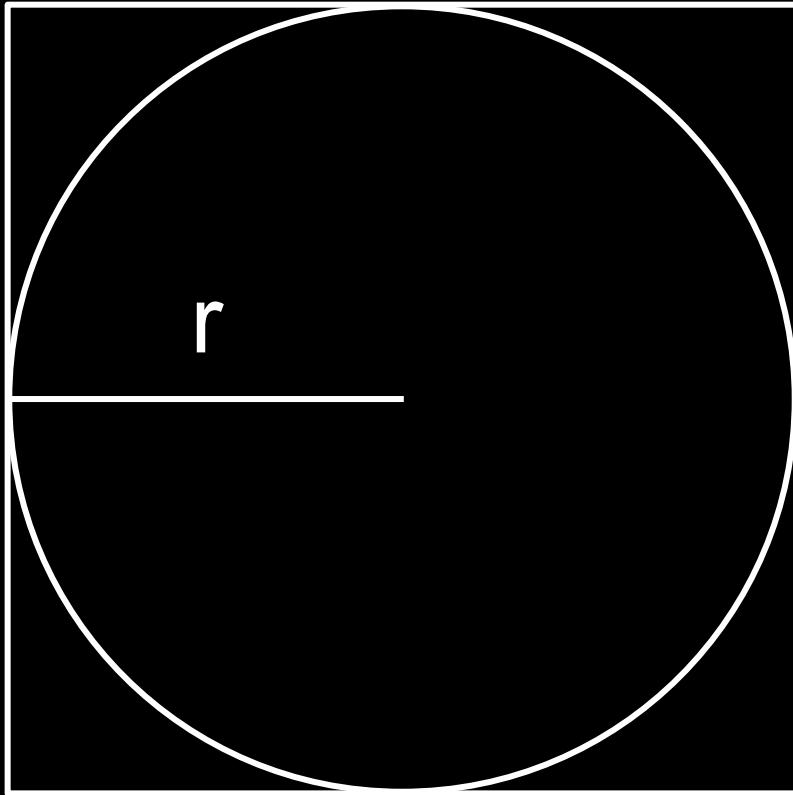
A diagram of a circle with a radius labeled  $r$ . The circle is drawn with a thin white line on a black background. A horizontal line segment from the center to the circumference represents the radius, with the letter  $r$  positioned at its midpoint.



 $2r$  $2r$

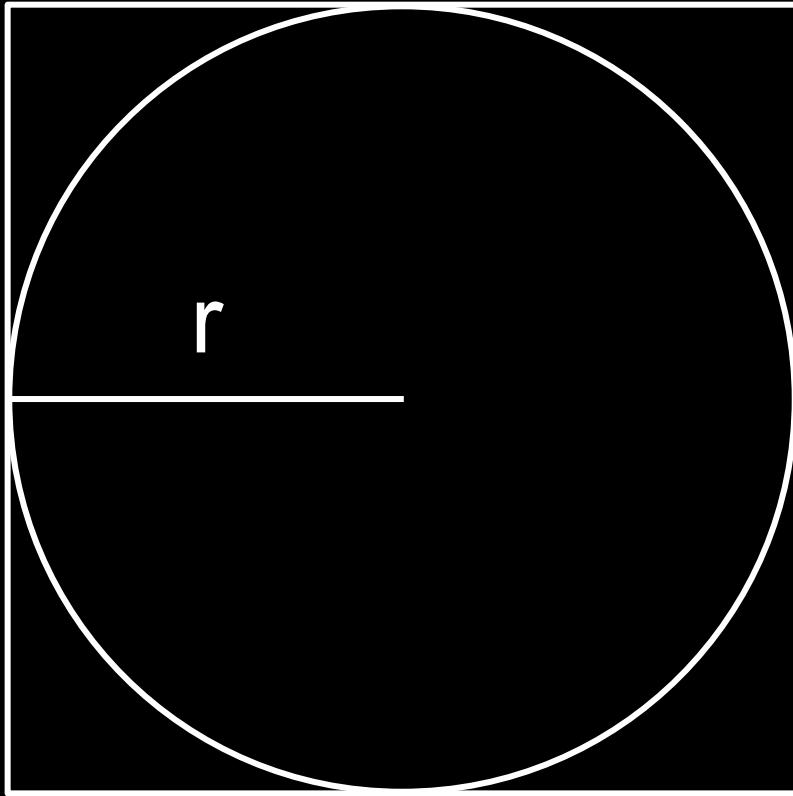


$$\frac{\textcircled{1}}{\square} = \frac{\pi r^2}{4r^2}$$



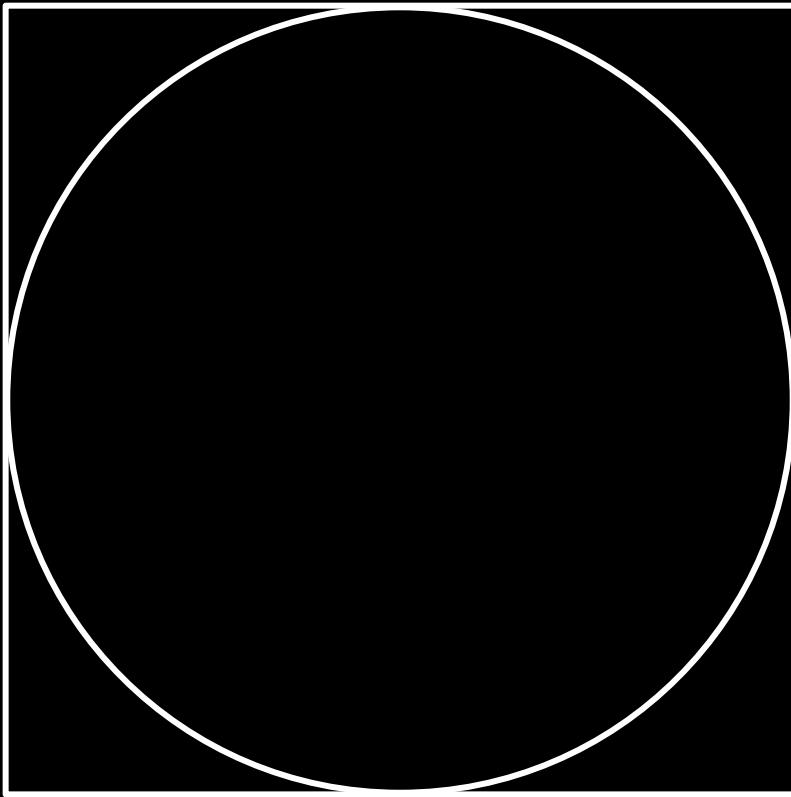
$2r$

$$\frac{\textcircled{1}}{\square} = \frac{\pi r^2}{4r^2}$$

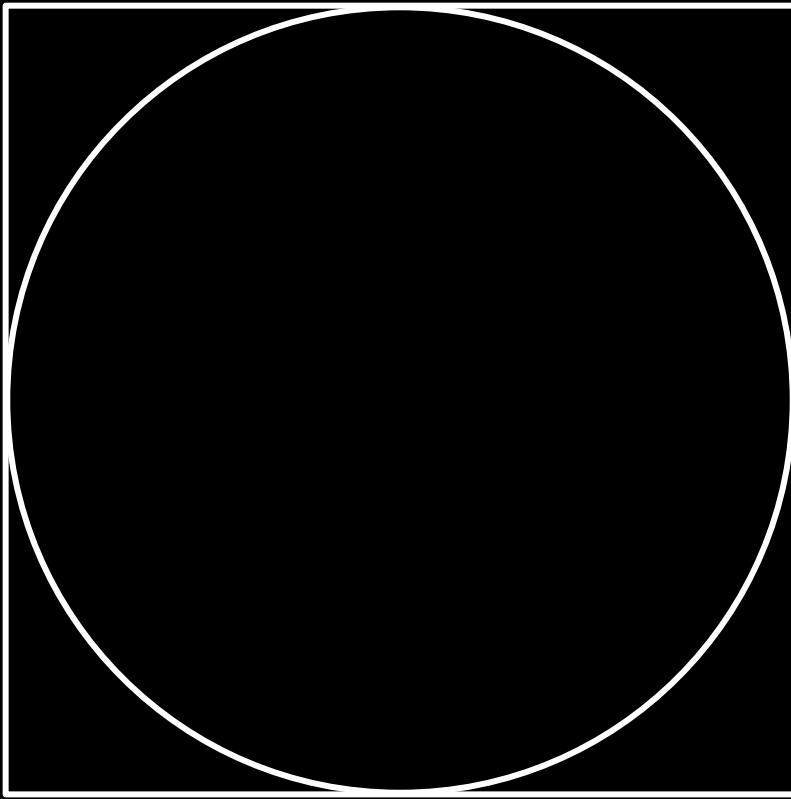


$$\frac{\textcircled{1}}{\square} = \frac{\pi}{4}$$

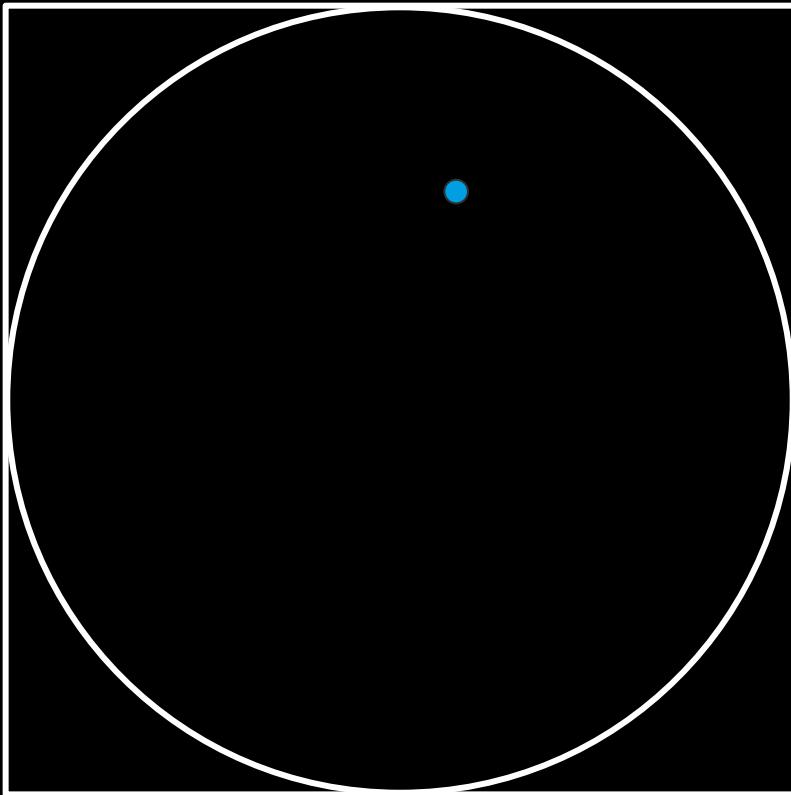
monte carlo simulation



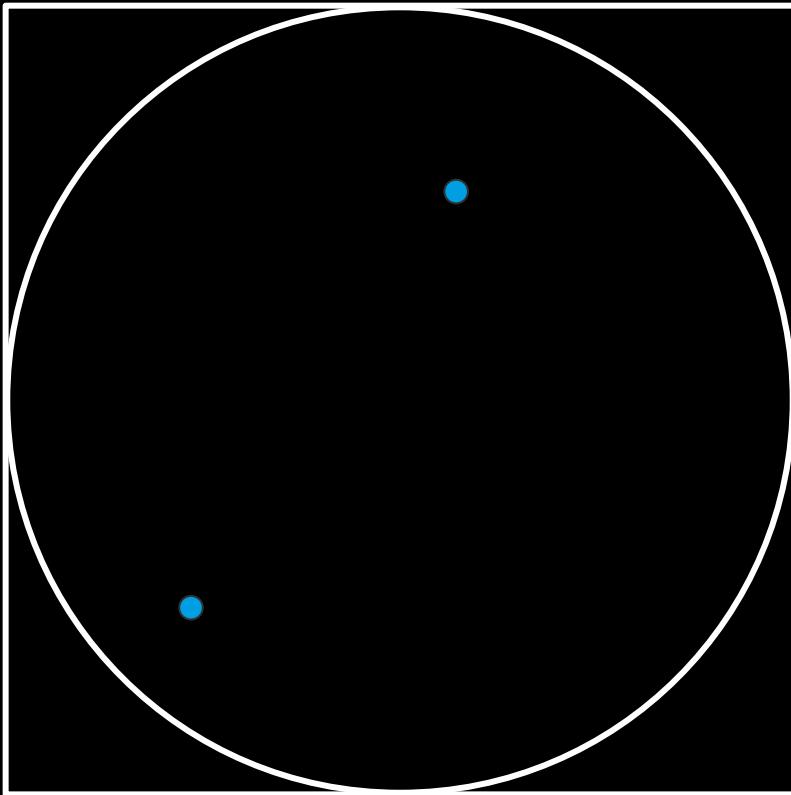
$$\frac{\textcircled{1}}{\square} = \frac{\pi}{4}$$



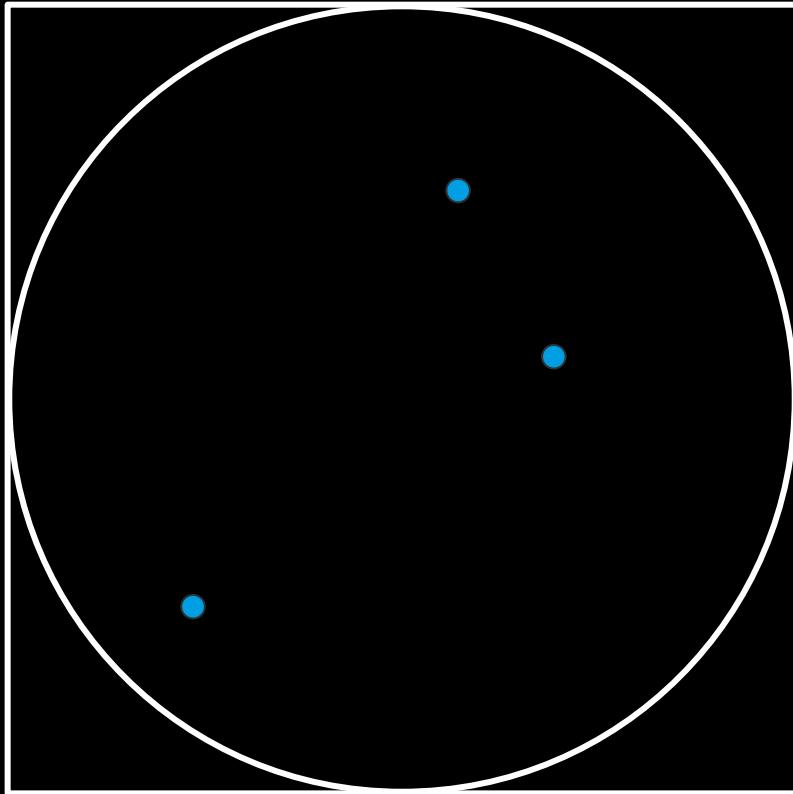
$$4 \frac{\bigcirc}{\square} = \pi$$



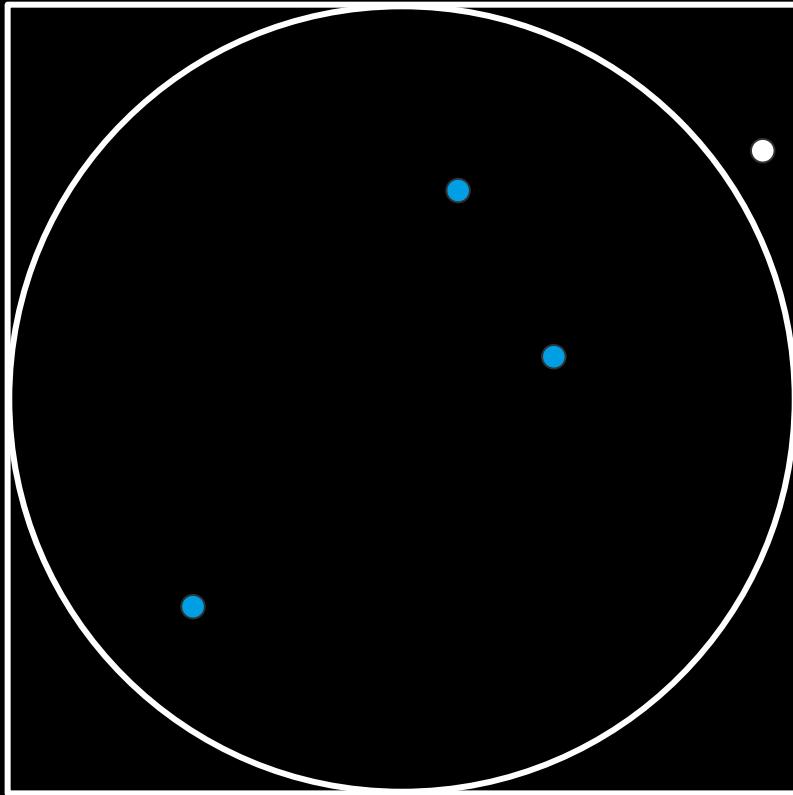
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 4$$



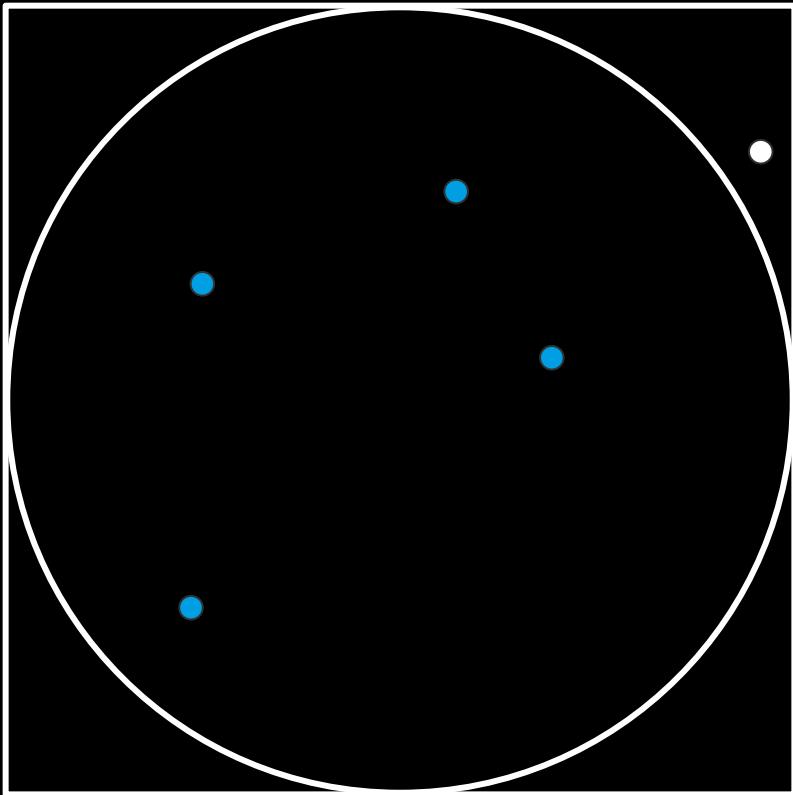
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 4$$



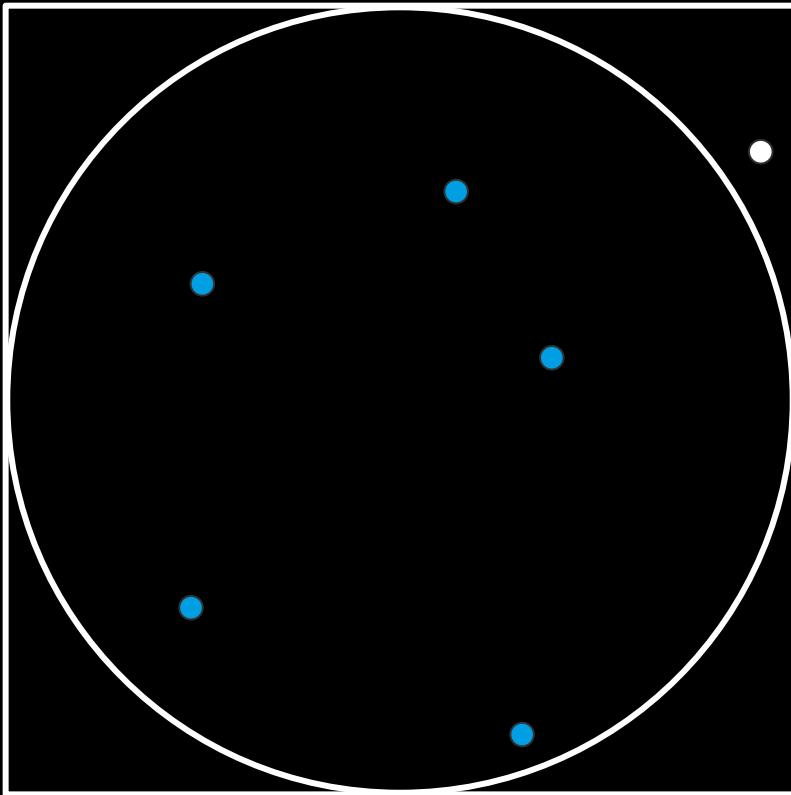
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 4$$



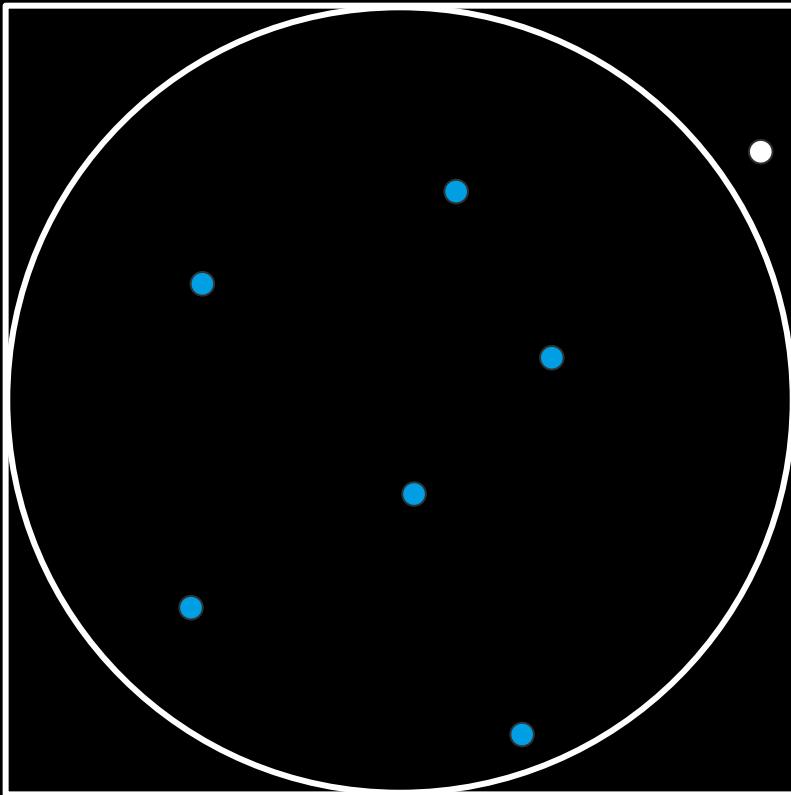
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 3$$



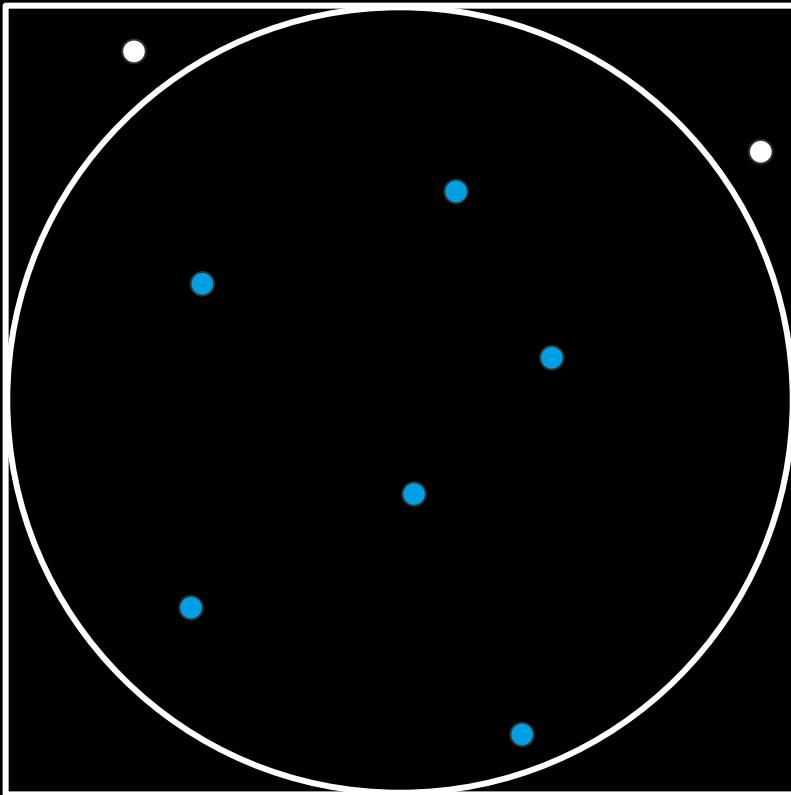
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 3,2$$



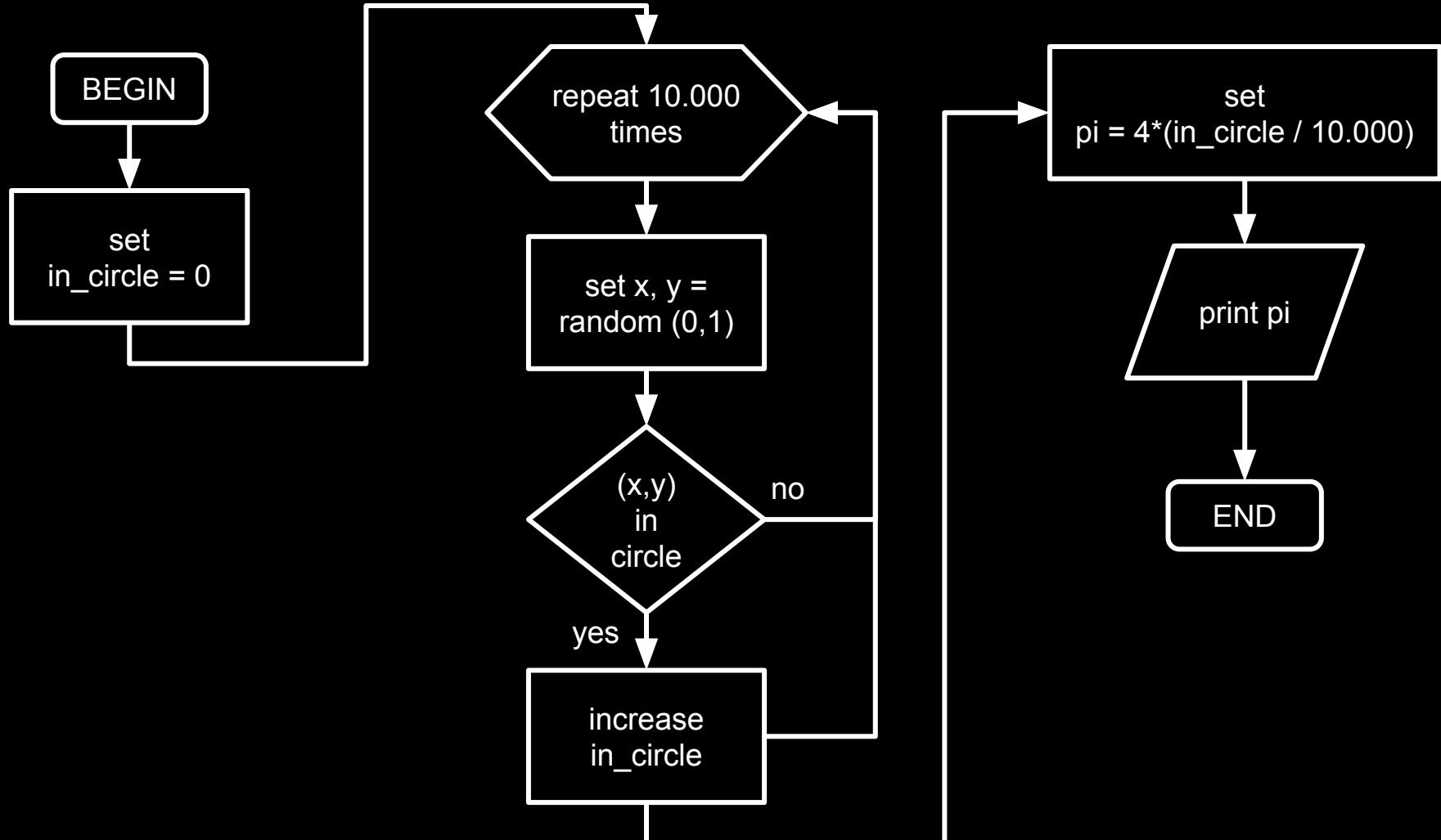
$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 3,33$$



$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 3,43$$



$$4 \frac{\bigcirc}{\square} = \pi$$
$$= 3$$



gregory-leibniz series

$$\pi = 4 \sum_{n=0}^{\infty} \frac{(-1)^n}{2n+1}$$

$$\pi = 4 \left( 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \frac{1}{9} - \frac{1}{11} + \dots \right)$$



# sorting

[ 9, 5, 2, 1, 4, 7 ]

bubble sort

repeatedly compare and swap elements until done.

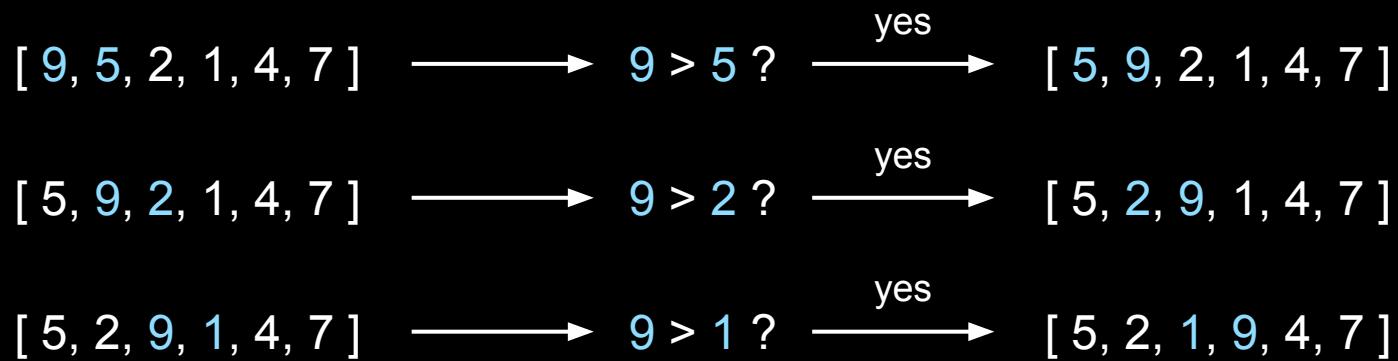
[ 9, 5, 2, 1, 4, 7 ]

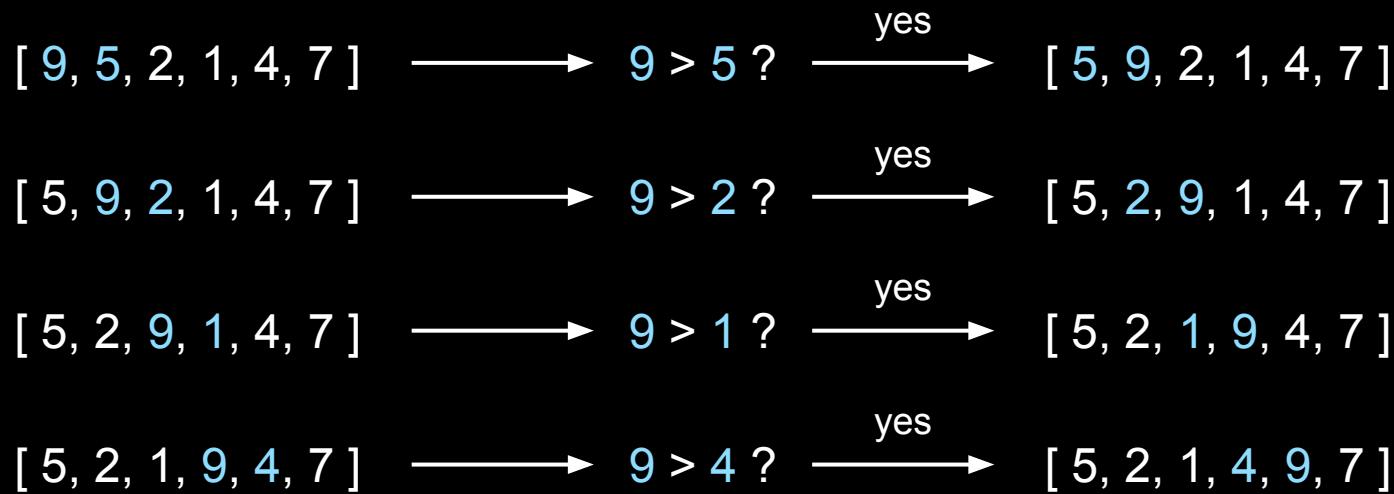
[ 9, 5, 2, 1, 4, 7 ]     $\longrightarrow$      $9 > 5 ?$

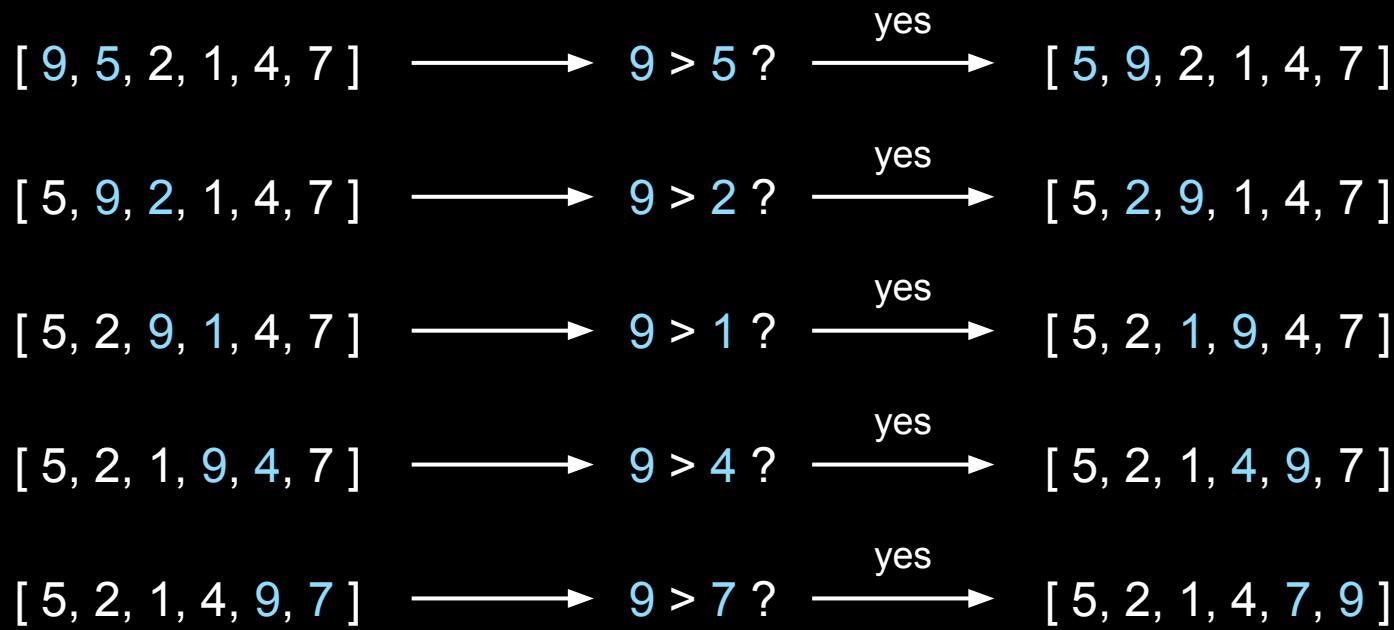
[ 9, 5, 2, 1, 4, 7 ] → 9 > 5 ? → yes [ 5, 9, 2, 1, 4, 7 ]

[ 9, 5, 2, 1, 4, 7 ] → 9 > 5 ? →  
yes [ 5, 9, 2, 1, 4, 7 ]

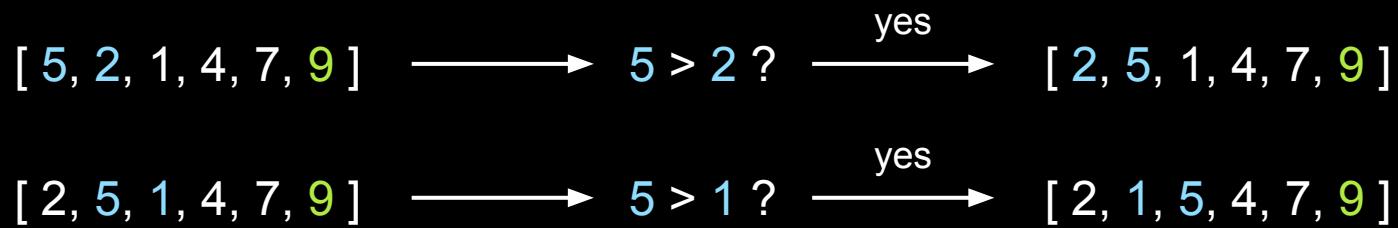
[ 5, 9, 2, 1, 4, 7 ] → 9 > 2 ? →  
yes [ 5, 2, 9, 1, 4, 7 ]

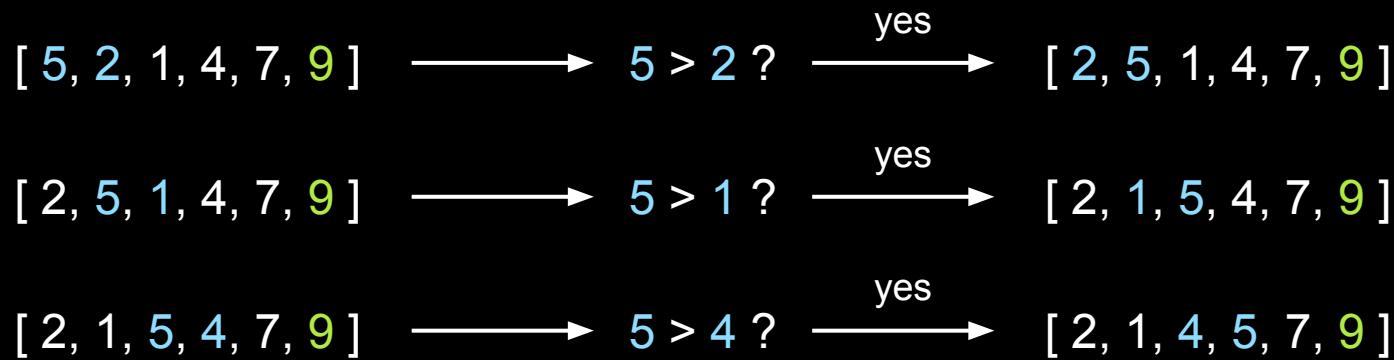


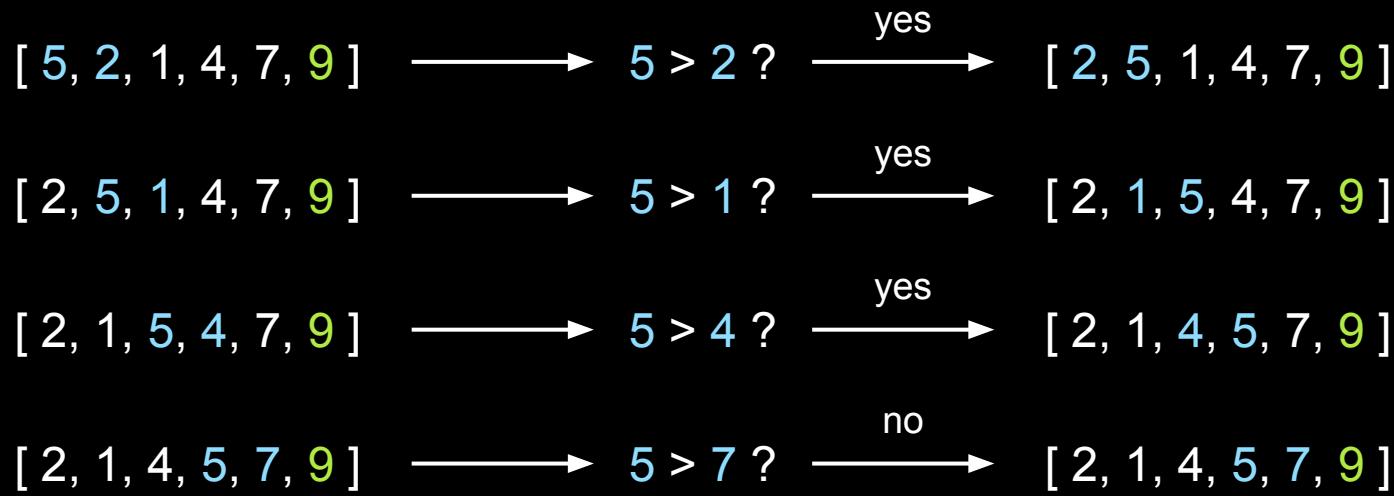




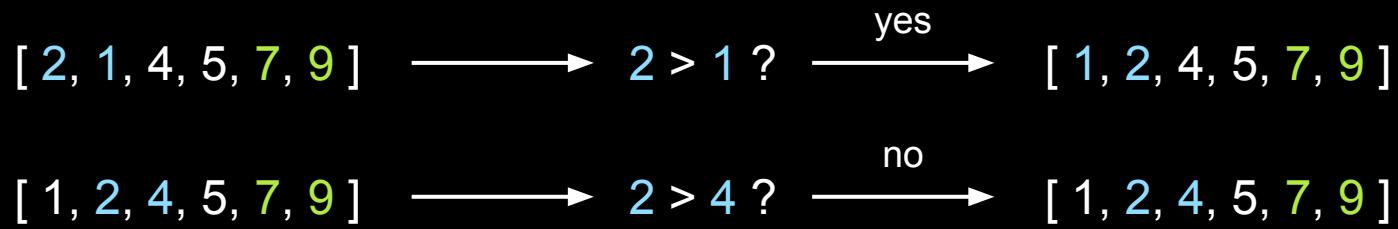
[ 5, 2, 1, 4, 7, 9 ] → 5 > 2 ? → yes [ 2, 5, 1, 4, 7, 9 ]

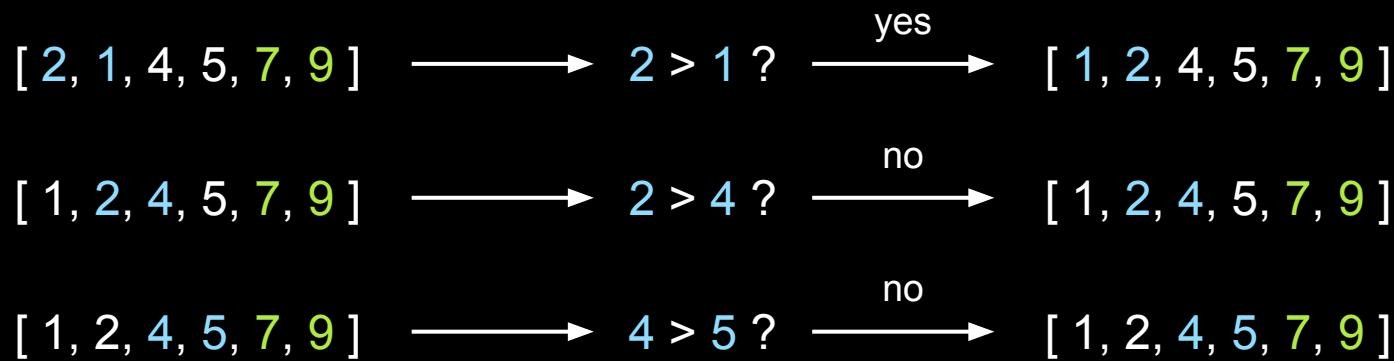






[ 2, 1, 4, 5, 7, 9 ] → 2 > 1 ? → yes [ 1, 2, 4, 5, 7, 9 ]





[ 1, 2, 4, 5, 7, 9 ] → 1 > 2 ? →  
no [ 1, 2, 4, 5, 7, 9 ]

[ 1, 2, 4, 5, 7, 9 ]  $\longrightarrow$  1 > 2 ?  $\xrightarrow{\text{no}}$  [ 1, 2, 4, 5, 7, 9 ]  
[ 1, 2, 4, 5, 7, 9 ]  $\longrightarrow$  2 > 4 ?  $\xrightarrow{\text{no}}$  [ 1, 2, 4, 5, 7, 9 ]

[ 1, 2, 4, 5, 7, 9 ] → 1 > 2 ? →  
no [ 1, 2, 4, 5, 7, 9 ]

[ 1, 2, 4, 5, 7, 9 ] → 1 > 2 ? →  
no [ 1, 2, 4, 5, 7, 9 ]

[ 1, 2, 4, 5, 7, 9 ] DONE!

# selection sort

find the smallest element and move it to front. repeat for the rest of the elements.

move to front

[ 9, 5, 2, 1, 4, 7 ] → [ 1, 5, 9, 2, 4, 7 ]

[ 9, 5, 2, 1, 4, 7 ] move to front [ 1, 5, 9, 2, 4, 7 ]

[ 1, 5, 9, 2, 4, 7 ] move to front [ 1, 2, 5, 9, 4, 7 ]

move to front

[ 9, 5, 2, 1, 4, 7 ] → [ 1, 5, 9, 2, 4, 7 ]

[ 1, 5, 9, 2, 4, 7 ] → [ 1, 2, 5, 9, 4, 7 ]

[ 1, 2, 5, 9, 4, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

move to front

[ 9, 5, 2, 1, 4, 7 ] → [ 1, 5, 9, 2, 4, 7 ]

[ 1, 5, 9, 2, 4, 7 ] → [ 1, 2, 5, 9, 4, 7 ]

[ 1, 2, 5, 9, 4, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

[ 1, 2, 4, 5, 9, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

move to front

[ 9, 5, 2, 1, 4, 7 ] → [ 1, 5, 9, 2, 4, 7 ]

[ 1, 5, 9, 2, 4, 7 ] → [ 1, 2, 5, 9, 4, 7 ]

[ 1, 2, 5, 9, 4, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

[ 1, 2, 4, 5, 9, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

[ 1, 2, 4, 5, 9, 7 ] → [ 1, 2, 4, 5, 7, 9 ]

move to front

[ 9, 5, 2, 1, 4, 7 ] → [ 1, 5, 9, 2, 4, 7 ]

[ 1, 5, 9, 2, 4, 7 ] → [ 1, 2, 5, 9, 4, 7 ]

[ 1, 2, 5, 9, 4, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

[ 1, 2, 4, 5, 9, 7 ] → [ 1, 2, 4, 5, 9, 7 ]

[ 1, 2, 4, 5, 9, 7 ] → [ 1, 2, 4, 5, 7, 9 ]

[ 1, 2, 4, 5, 7, 9 ] → [ 1, 2, 4, 5, 7, 9 ]

merge sort

divide the elements recursively in two halves until only one element is left.  
then merge the sorted halves back together.

divide the elements **recursively** in two halves until only one element is left.  
then merge the sorted halves back together.

[ 9, 5, 2, 1, 4, 7 ]

[ 9, 5, 2 ]

split

[ 1, 4, 7 ]

[ 9, 5, 2, 1, 4, 7 ]

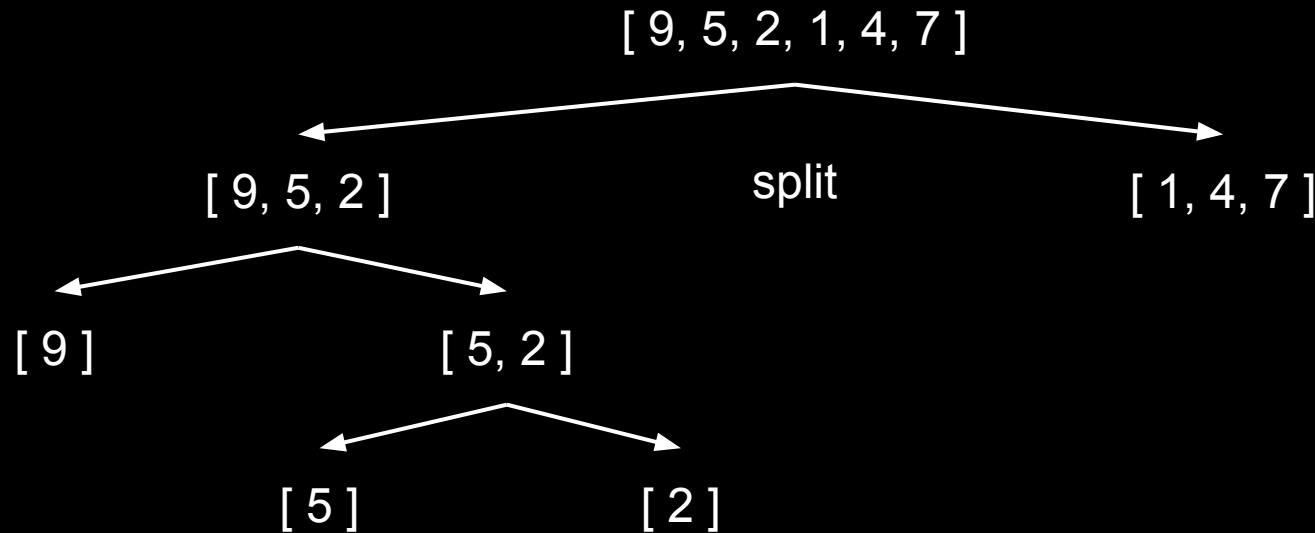
[ 9, 5, 2 ]

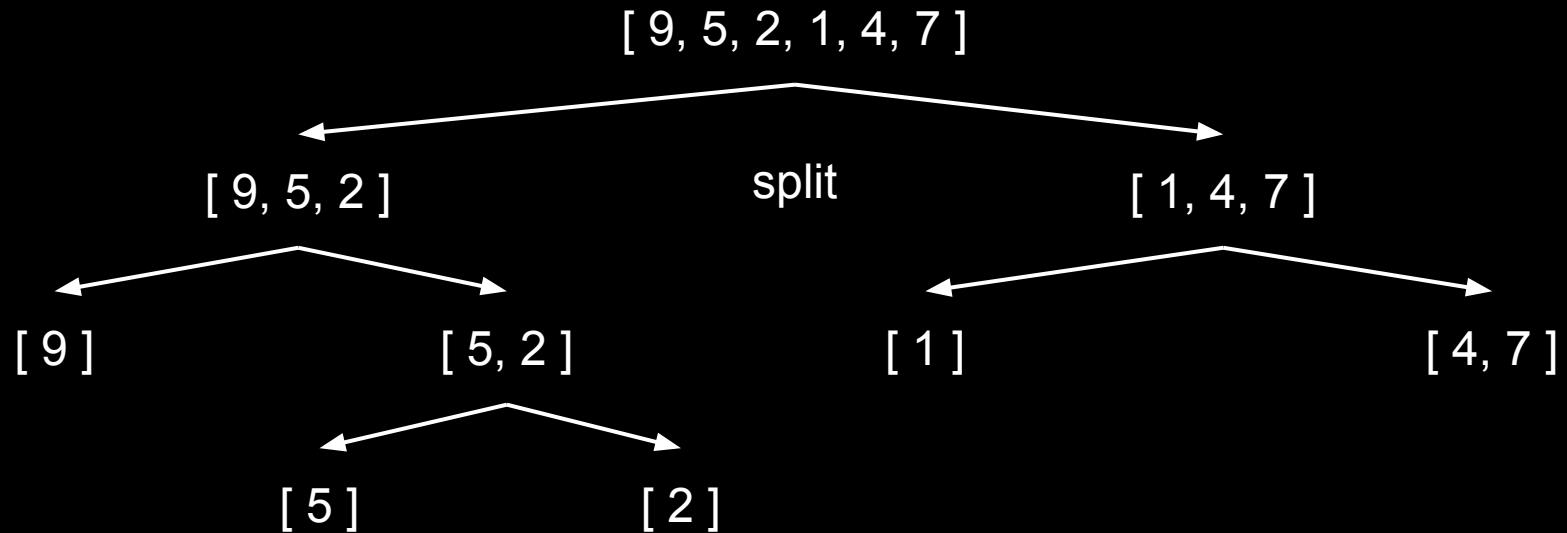
split

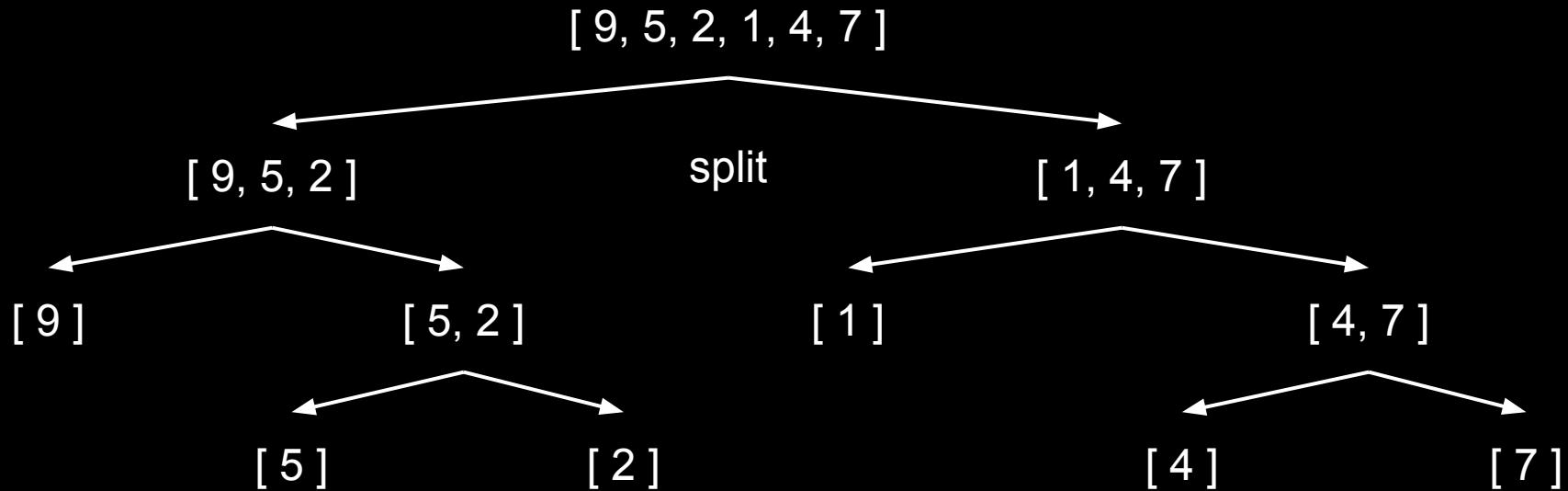
[ 1, 4, 7 ]

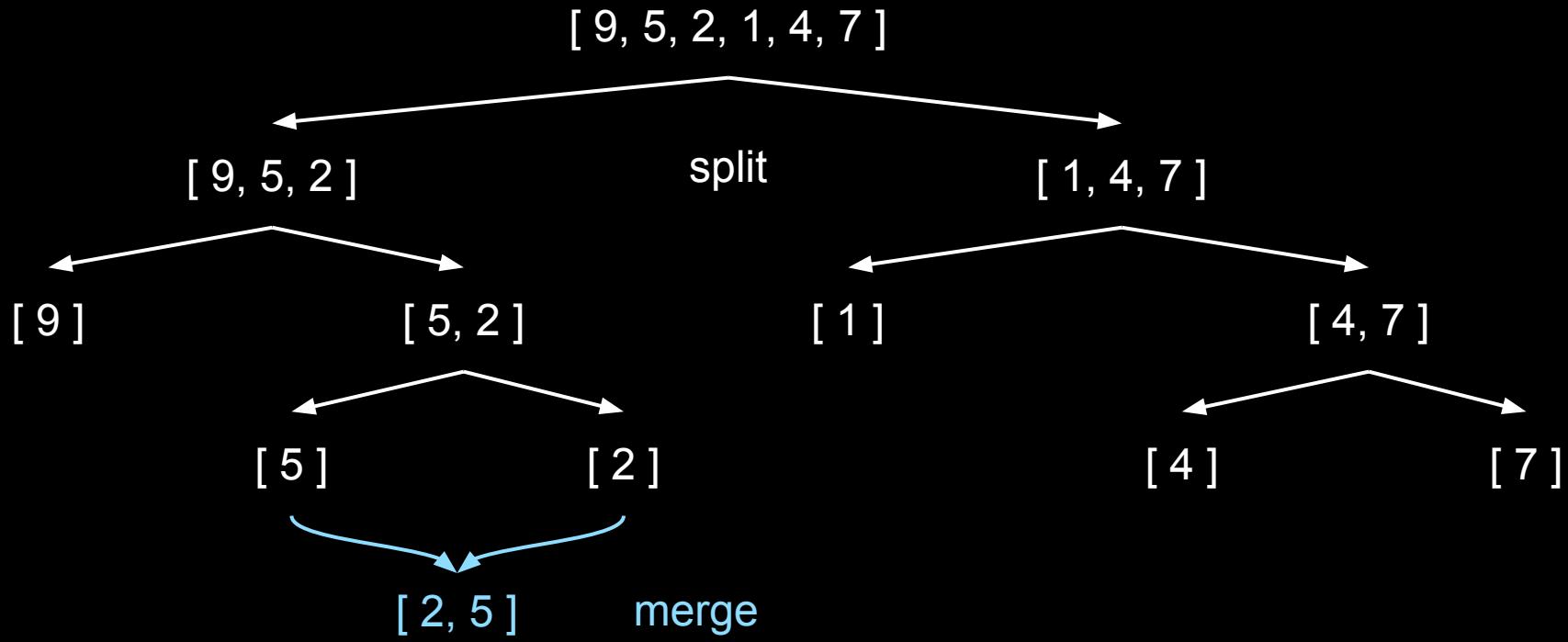
[ 9 ]

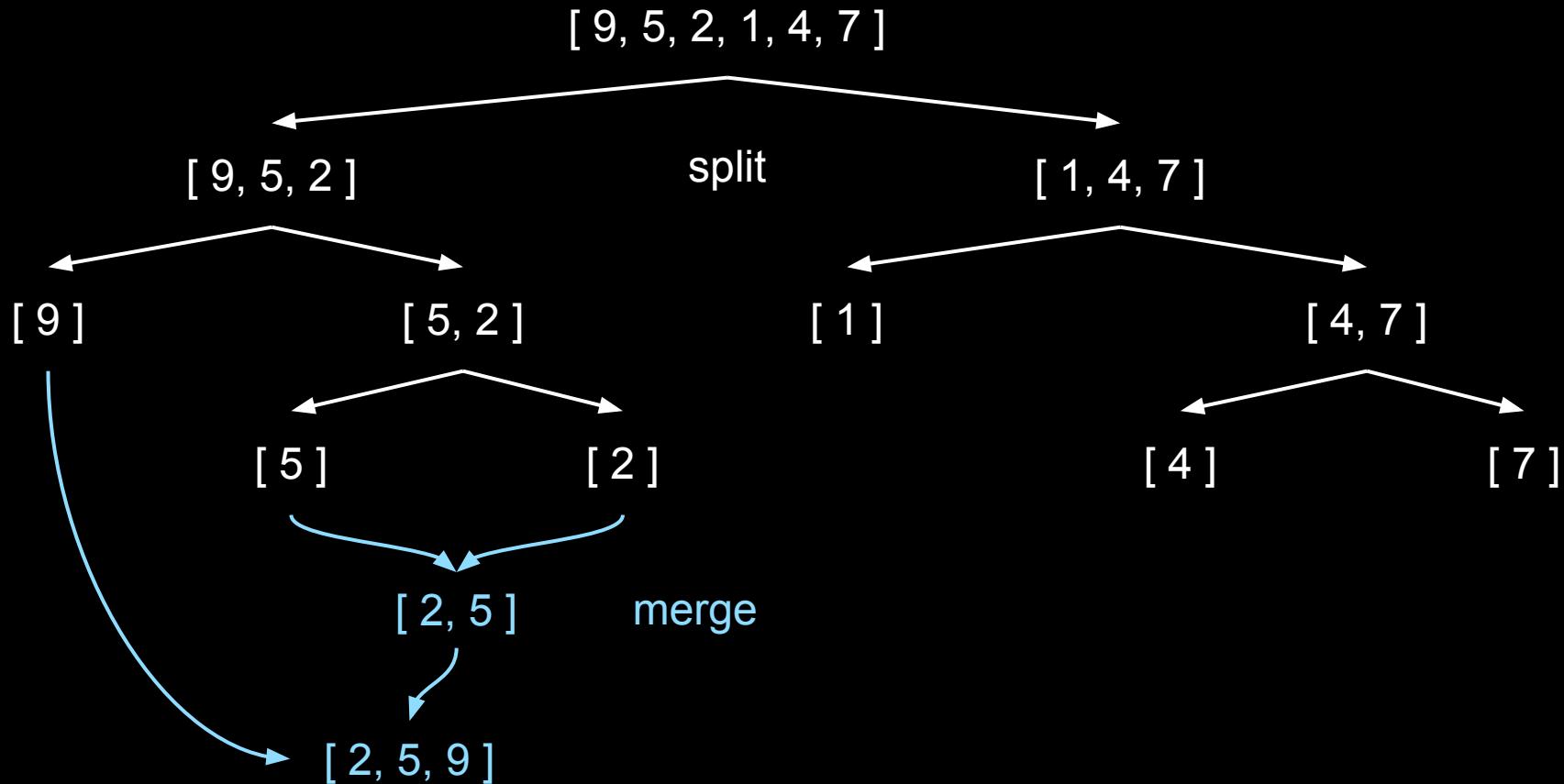
[ 5, 2 ]

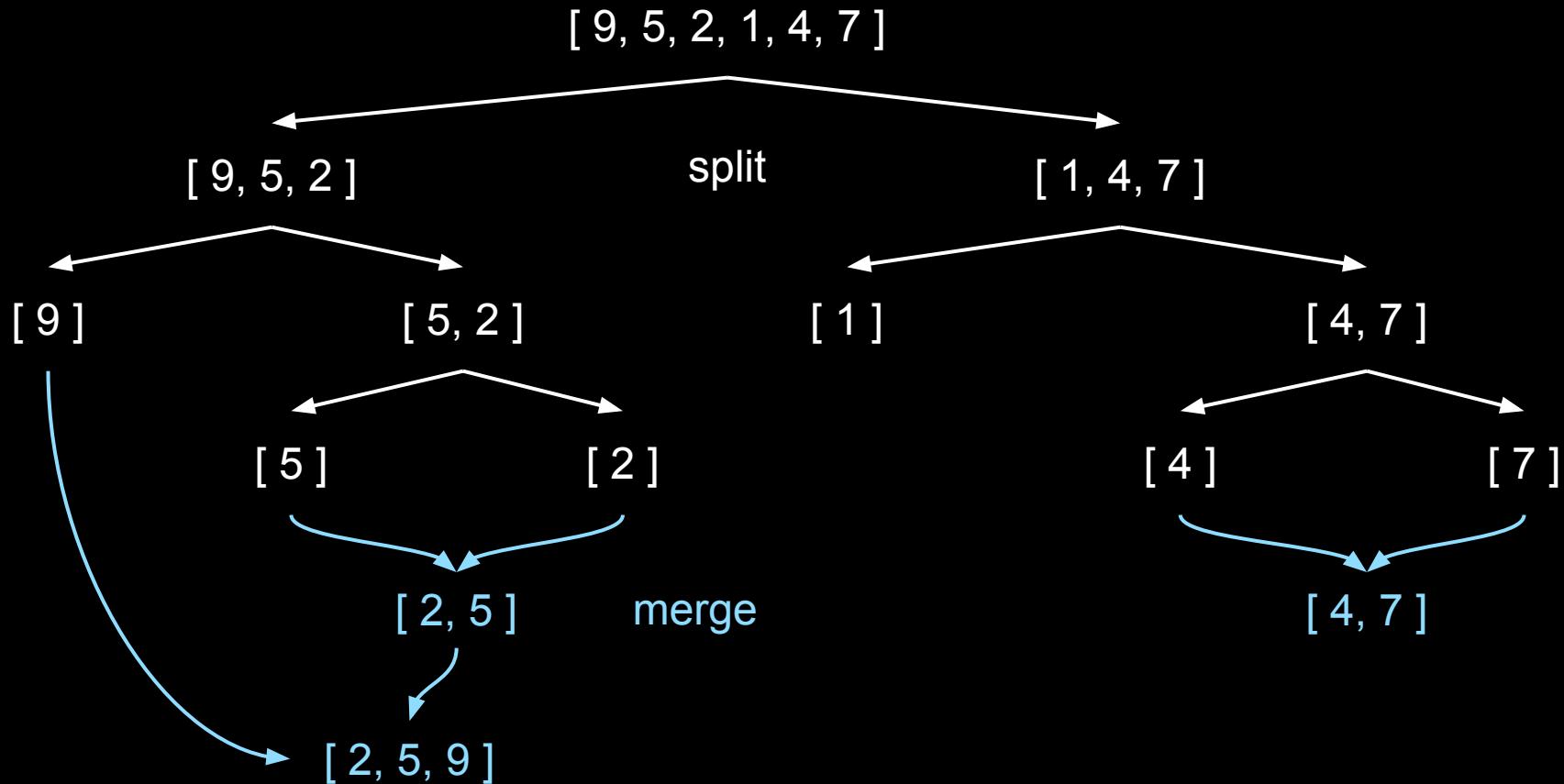


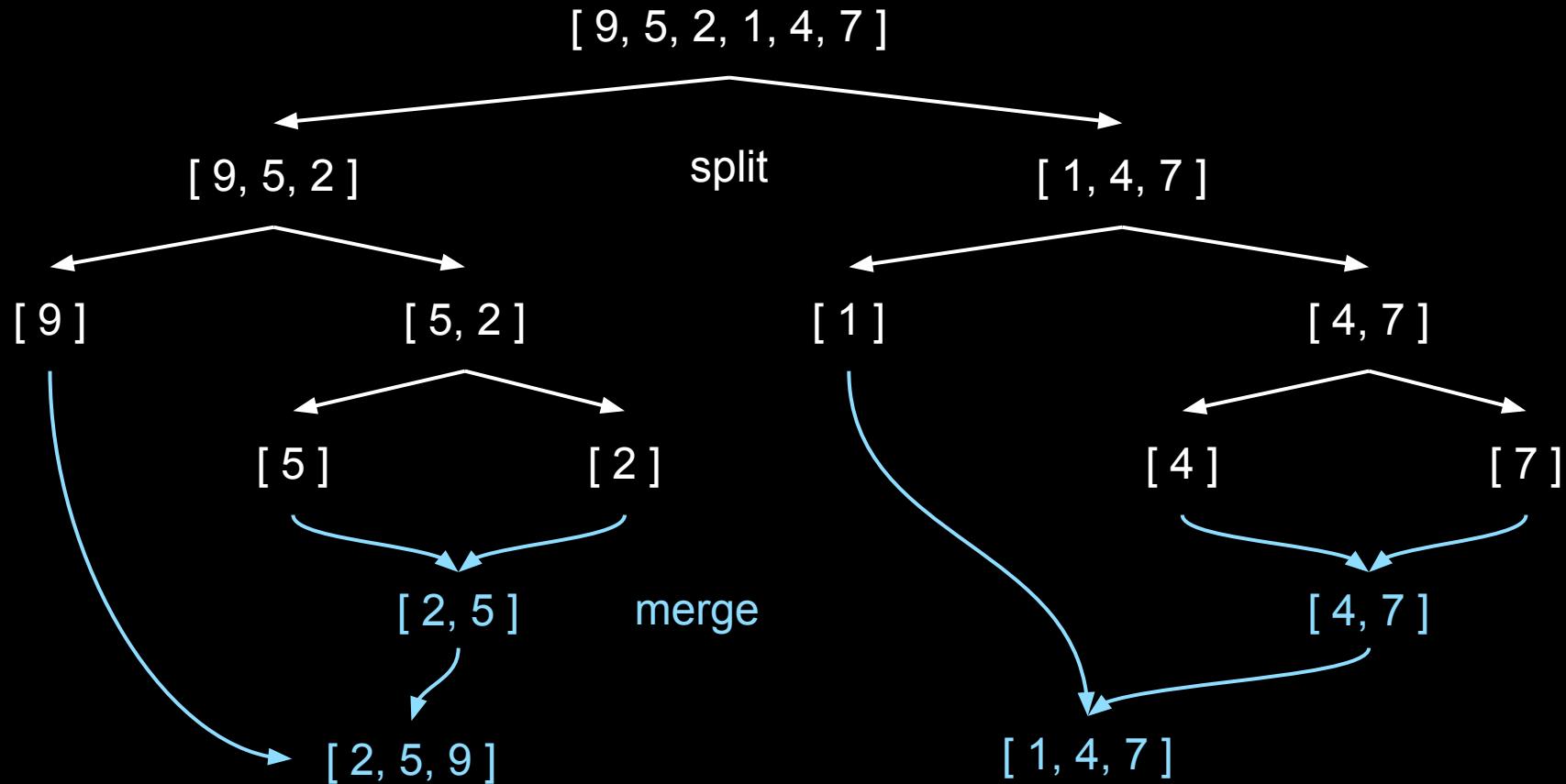


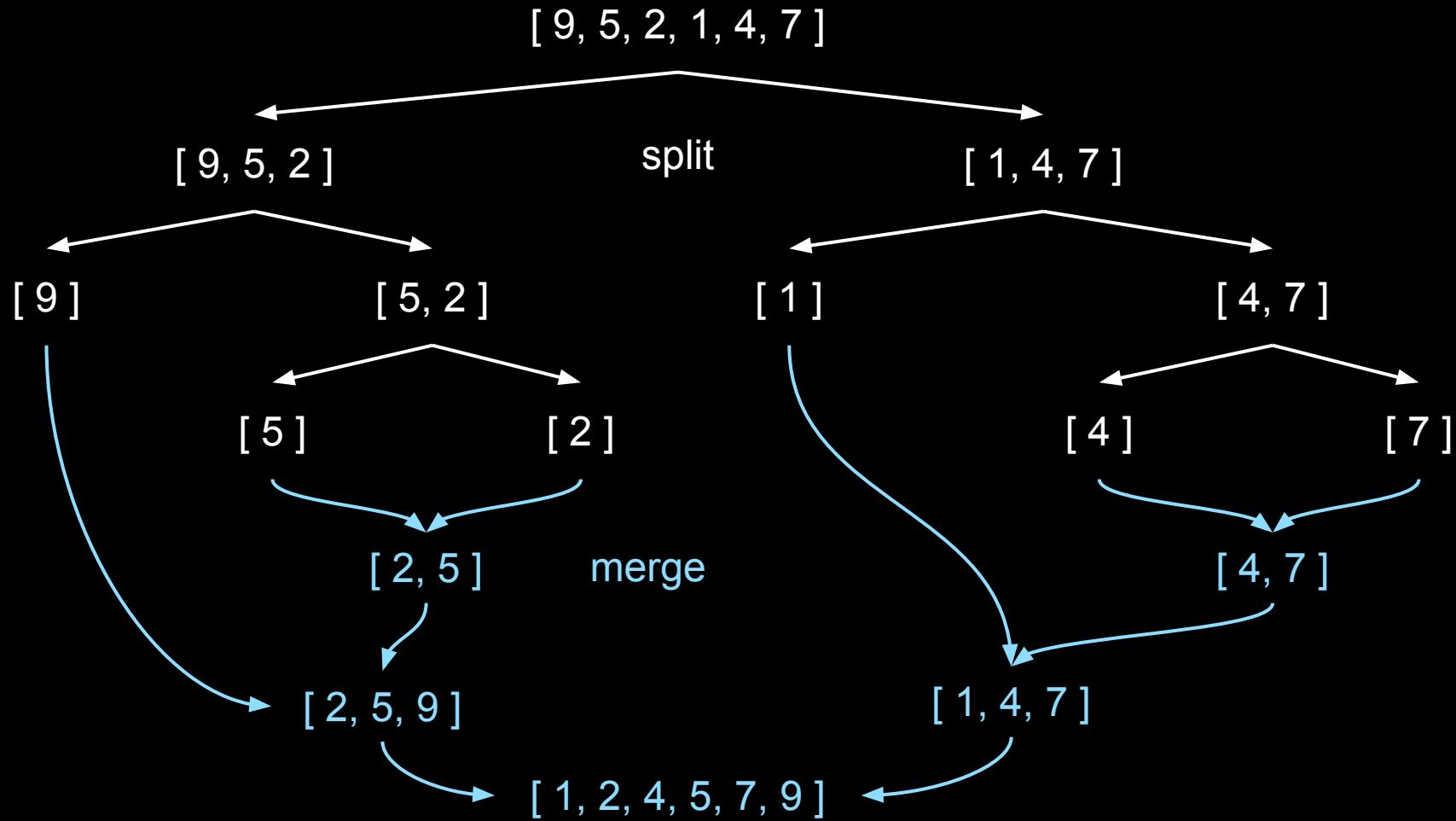




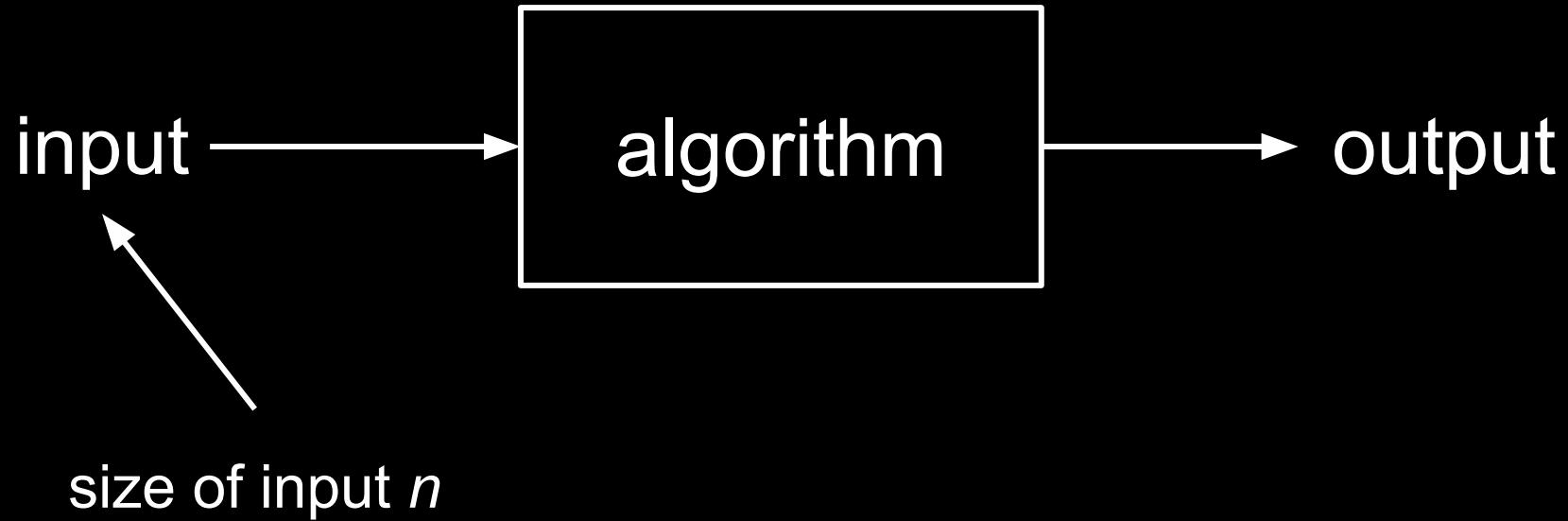


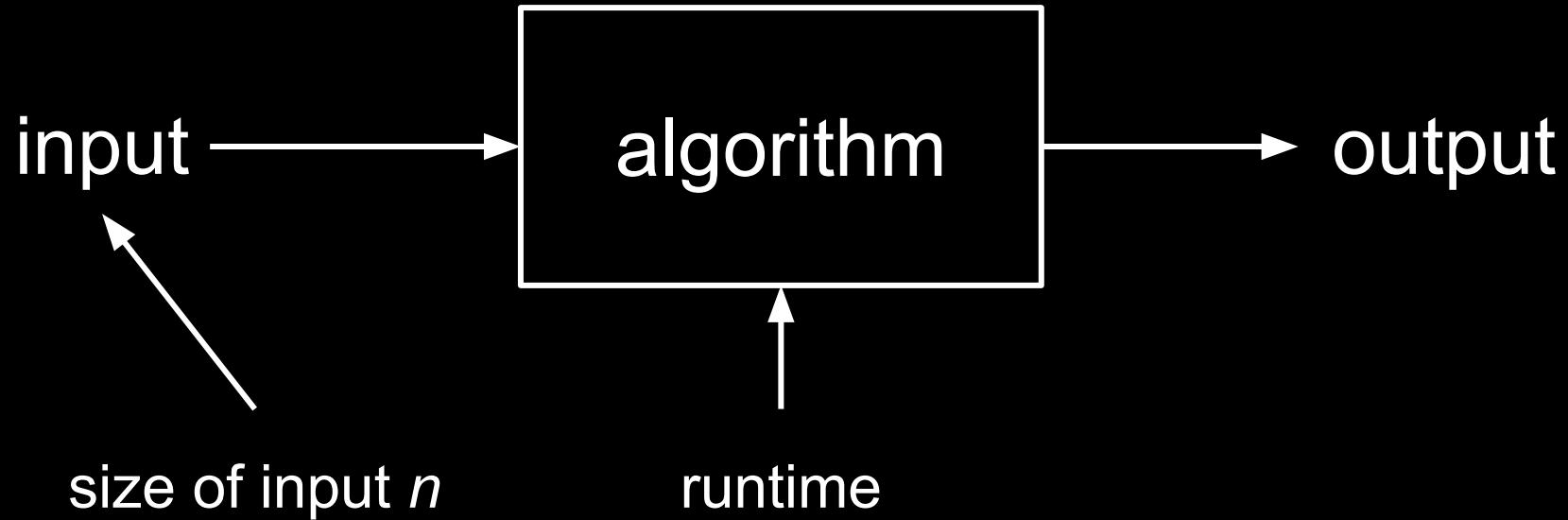




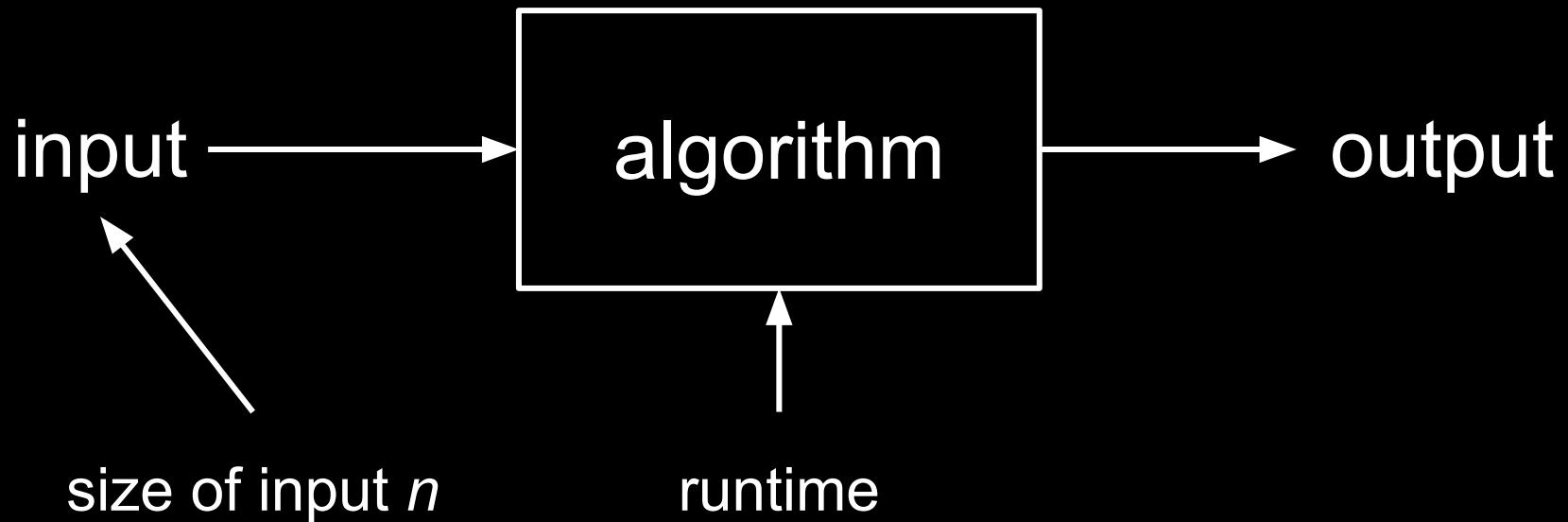


complexity





$O(n)$



$O(1)$	runtime is constant and independent of problem size
$O(\log_2 n)$	runtime is determined by the logarithm of problem size
$O(n)$	runtime is linear to problem size
$O(n^2)$	runtime grows quadratically with the size of the problem
$O(n^3)$	runtime grows cubically with the size of the problem
$O(2^n)$	runtime grows exponentially with the size of the problem
$O(n!)$	runtime grows factorially with the size of the problem

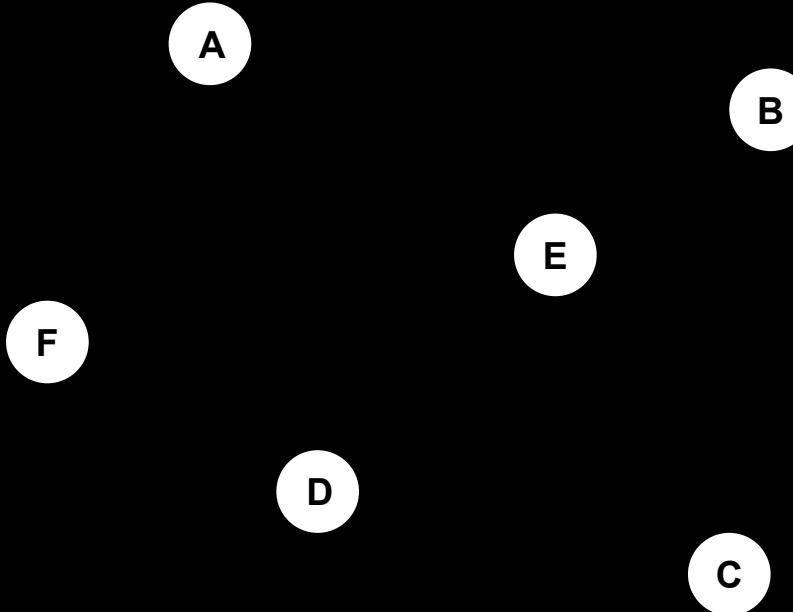


# optimization

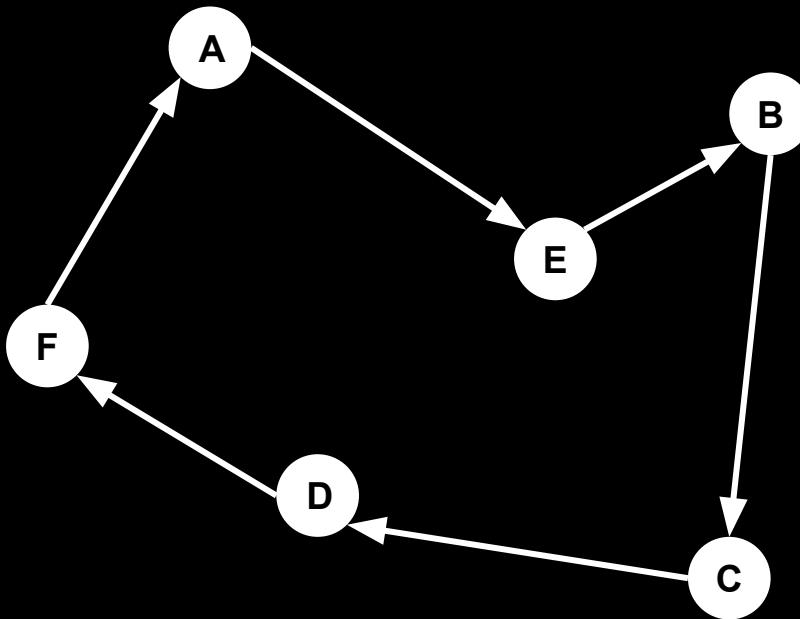


# traveling salesmen

# shortest tour?



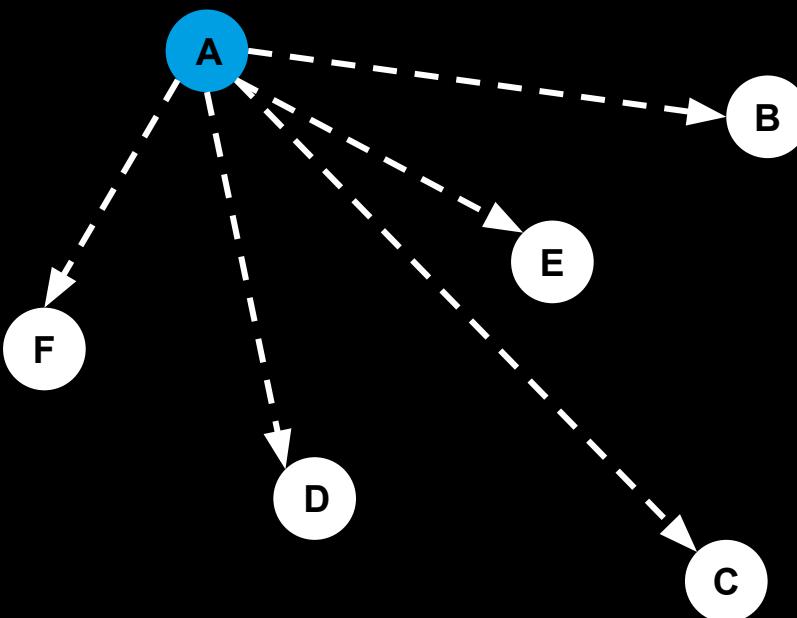
shortest tour?



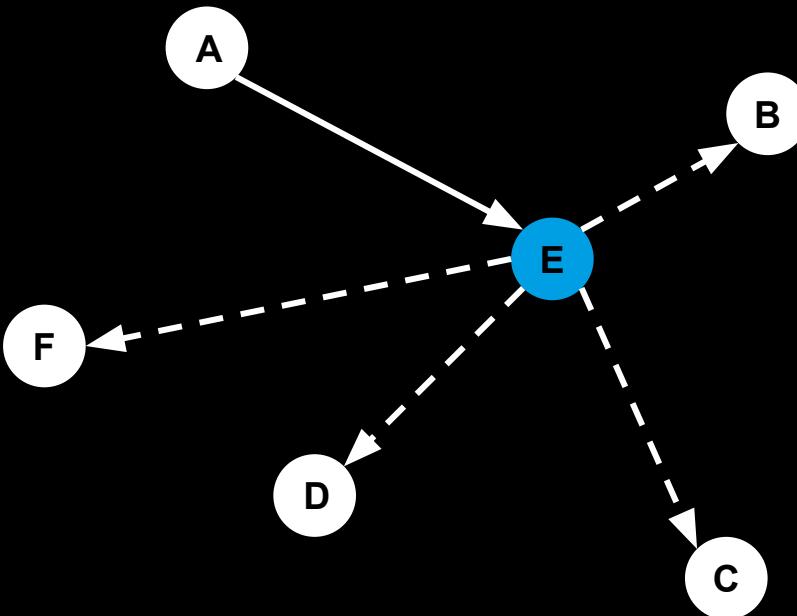
brute force

$O(n) = n!$

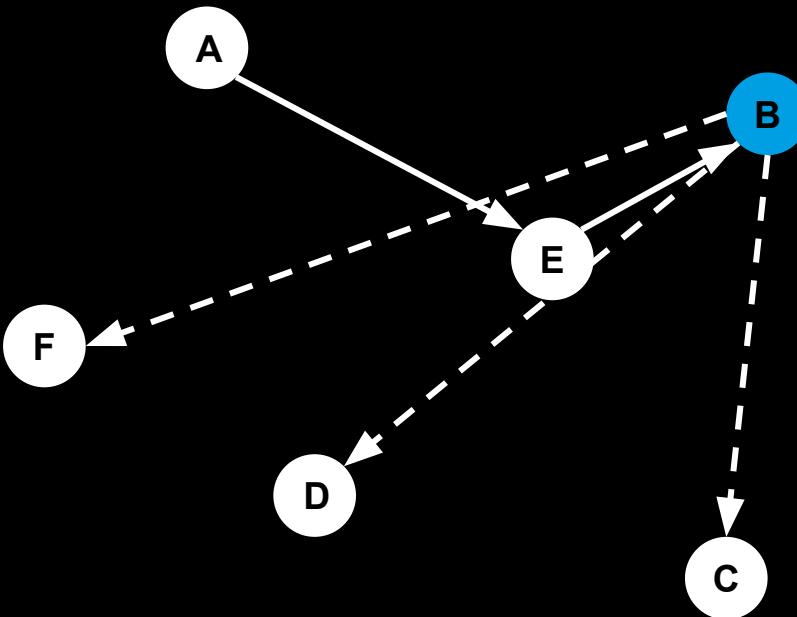
# 5 possible cities



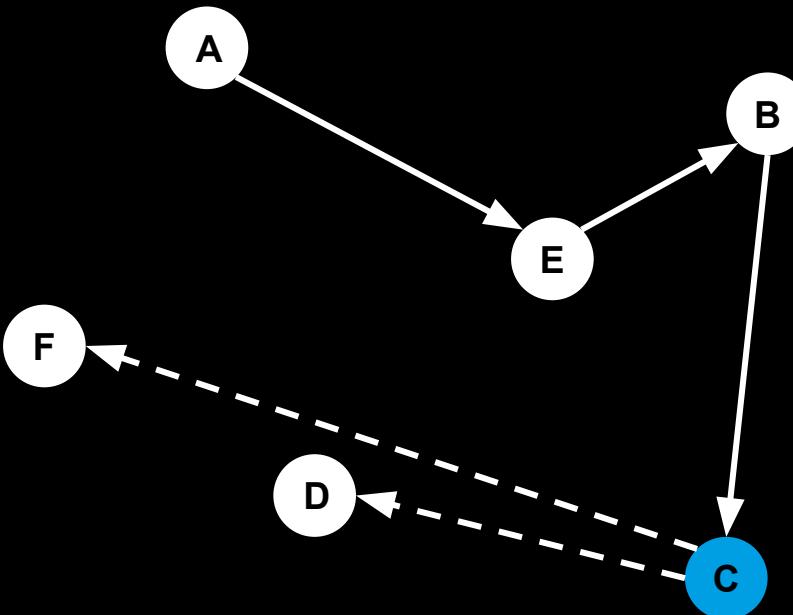
5 4 possible cities



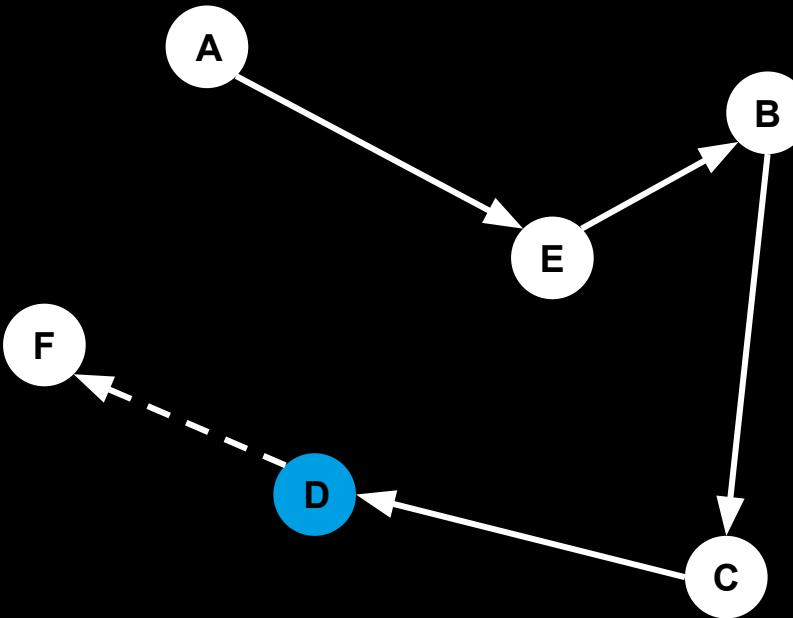
5 4 3 possible cities



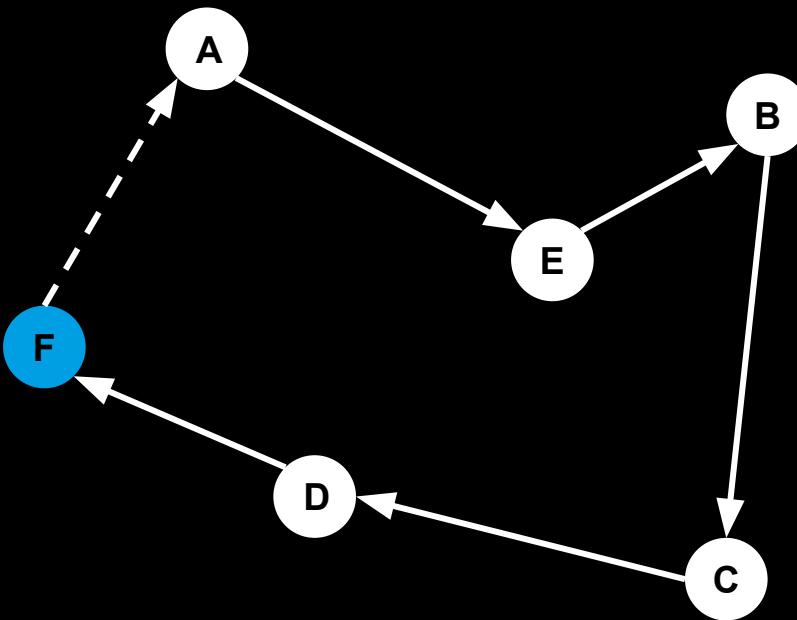
5 4 3 2 possible cities



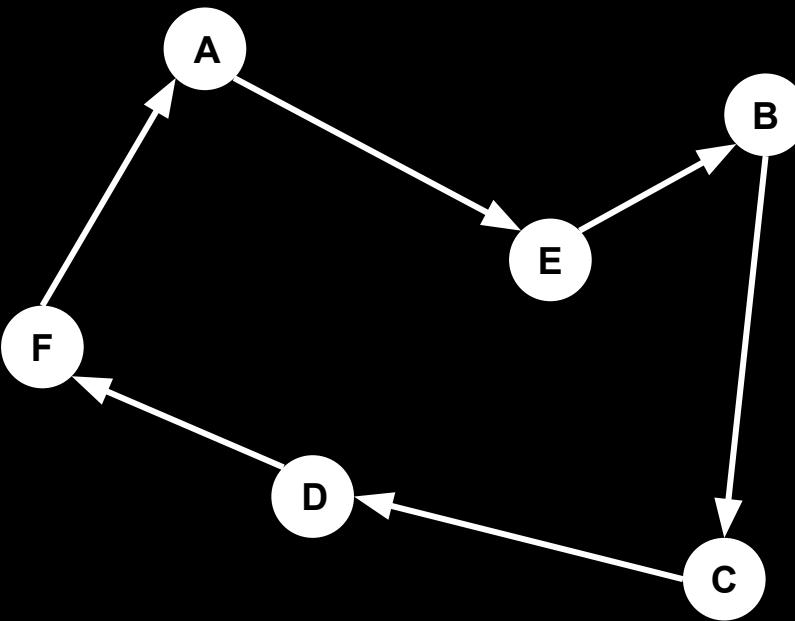
5 4 3 2 1 possible city



5 4 3 2 1 return nome



5 4 3 2 4 return nome



$$O(n) = n!$$

$$n = 5$$

$$O(n) = n!$$

$$n = 5$$

$$N = 5 * 4 * 3 * 2 * 1$$

$$O(n) = n!$$

$$n = 5$$

$$N = 5 * 4 * 3 * 2 * 1$$

$$= 120$$

$n = 10$

$n = 20$

$n = 30$

$$n = 25$$

brute force takes longer than the universe is old

$$n = 60$$

more possible routes than atoms in the universe

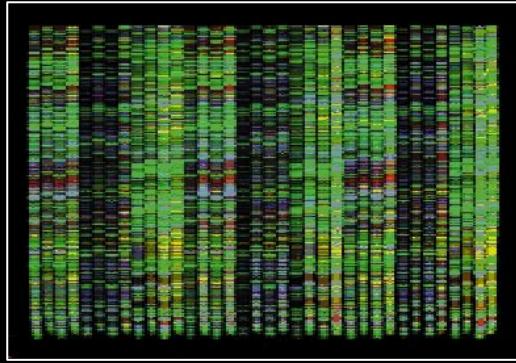


Image source: [IEEE](#)



Image source: [VDI Nachrichten](#)



Image source: [Wikimedia](#)



Image source: <https://github.com/meiyi1986/tutorials/blob/master/notebooks/img/pcb-drilling.jpeg>



Image source: [IAS Observatory](#)

but sometimes, brute force works perfectly



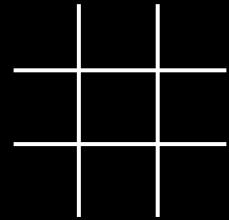
next\_move()

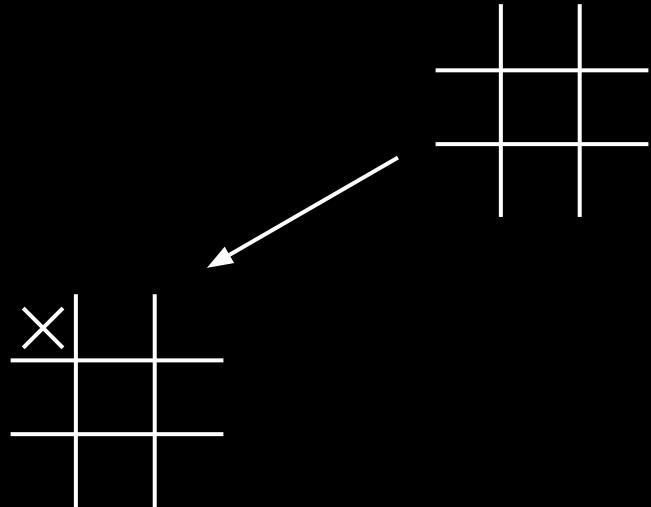
(2, 2)

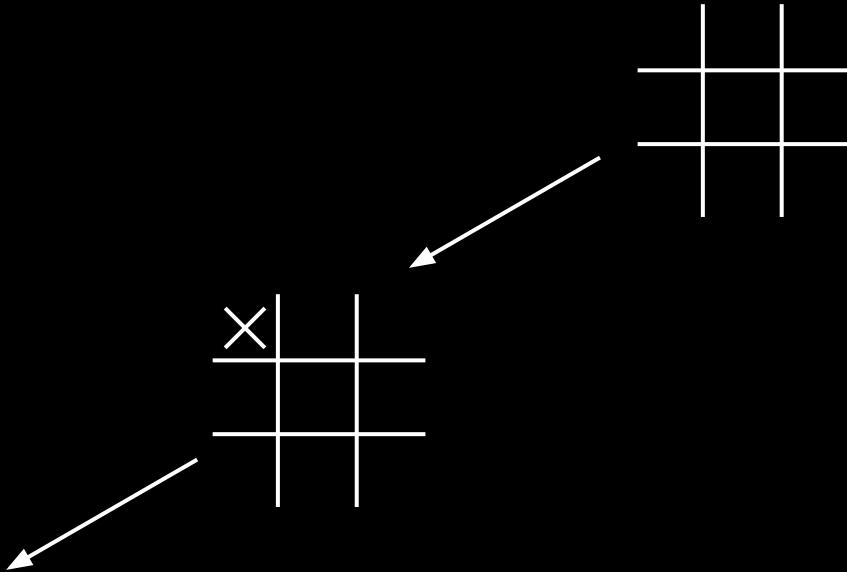
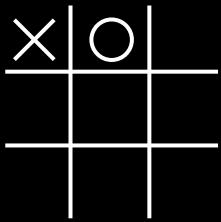
**brute-force search:**

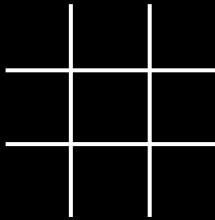
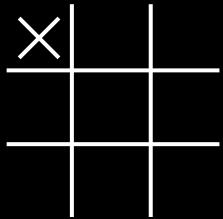
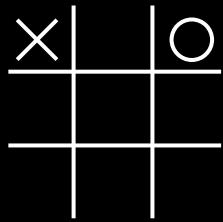
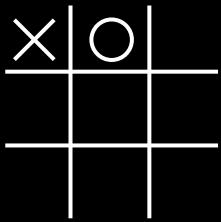
evaluate all possible move sequences.

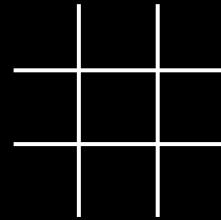
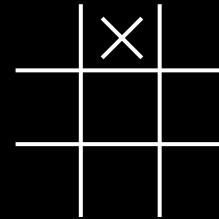
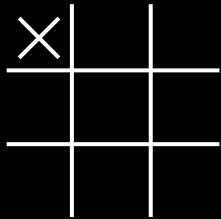
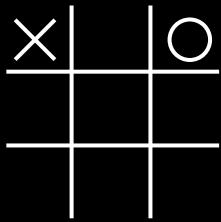
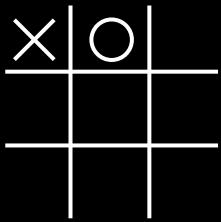
choose the move with the best value.

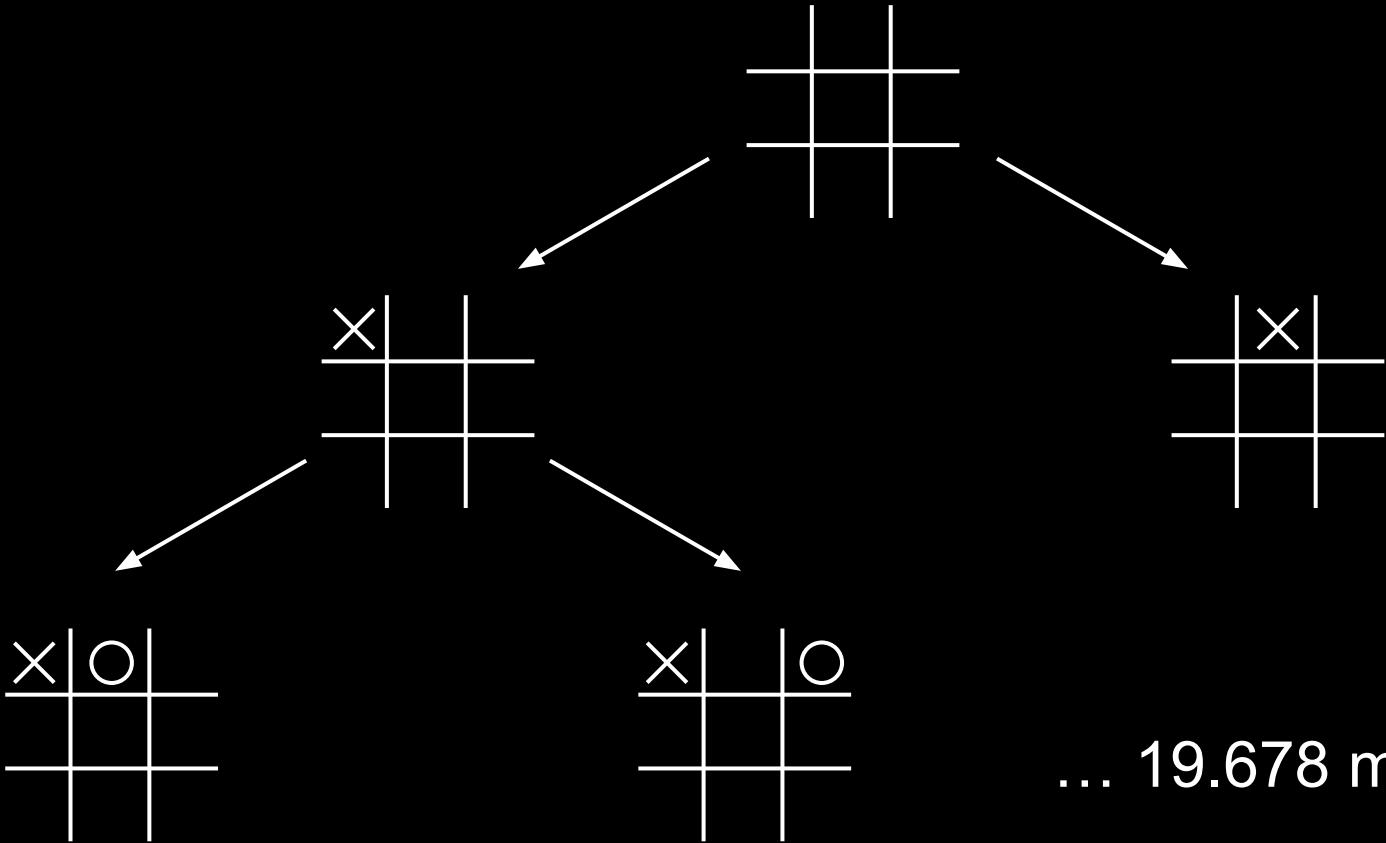








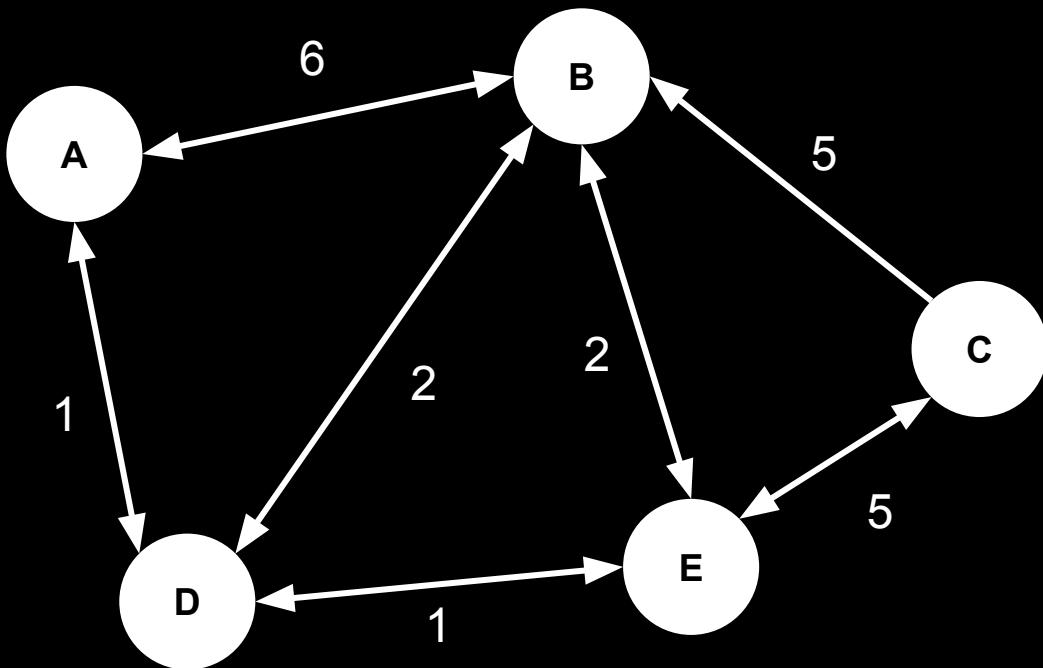




... 19.678 more



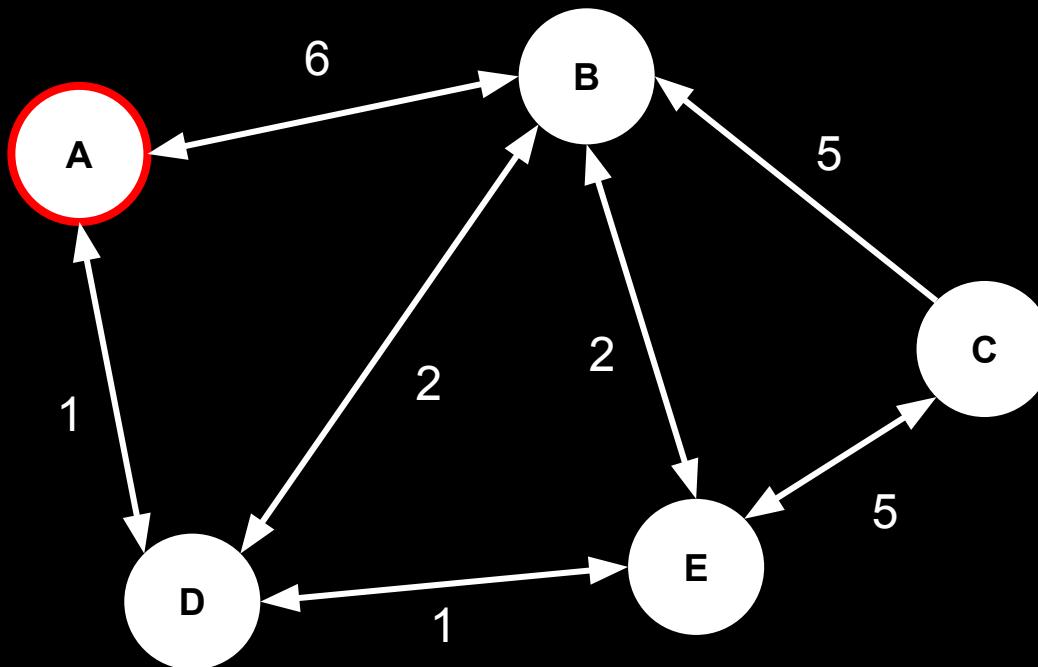
# shortest paths



# dijkstra's algorithm

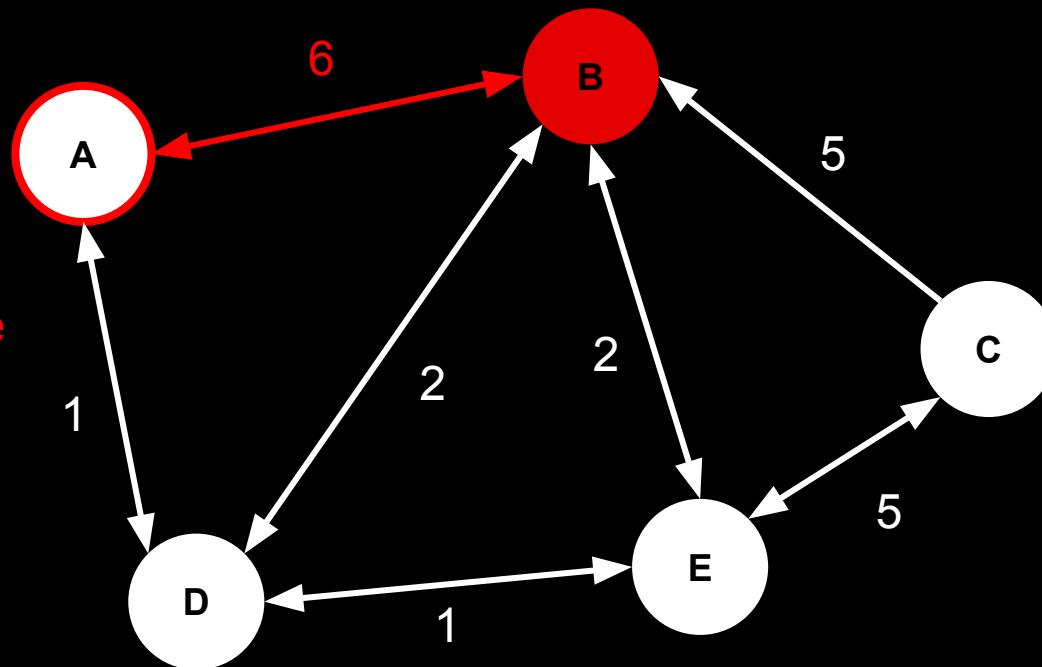
Distances: A = 0, B =  $\infty$ , C =  $\infty$ , D =  $\infty$ , E =  $\infty$

Set A as first location



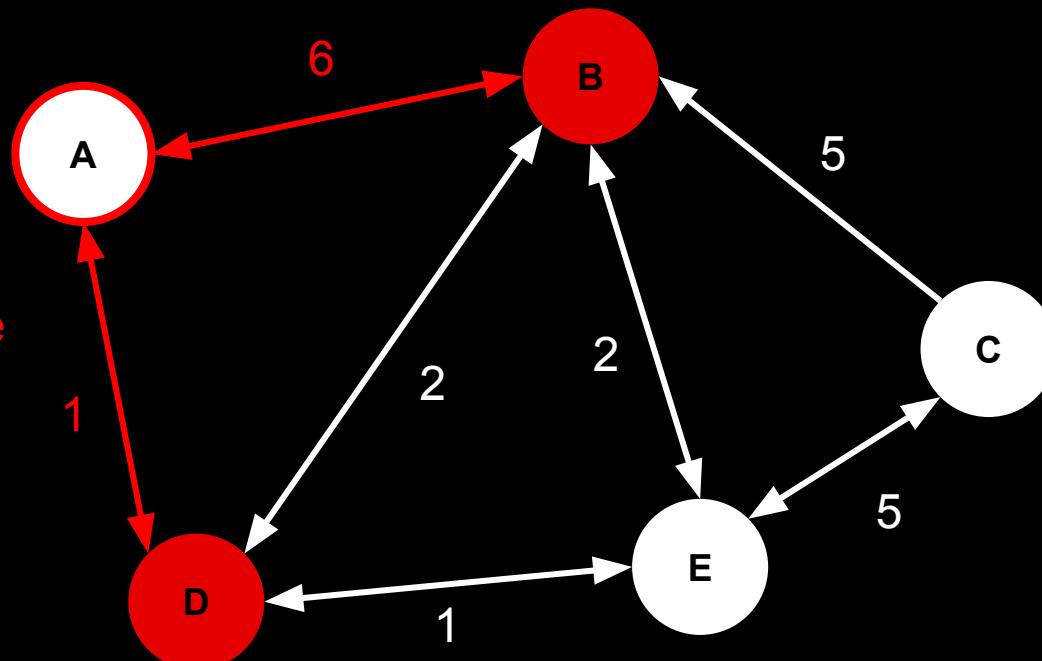
Distances: A = 0, B = 6, C =  $\infty$ , D =  $\infty$ , E =  $\infty$

Update distance  
to B to 6



Distances: A = 0, B = 6, C =  $\infty$ , D = 1, E =  $\infty$

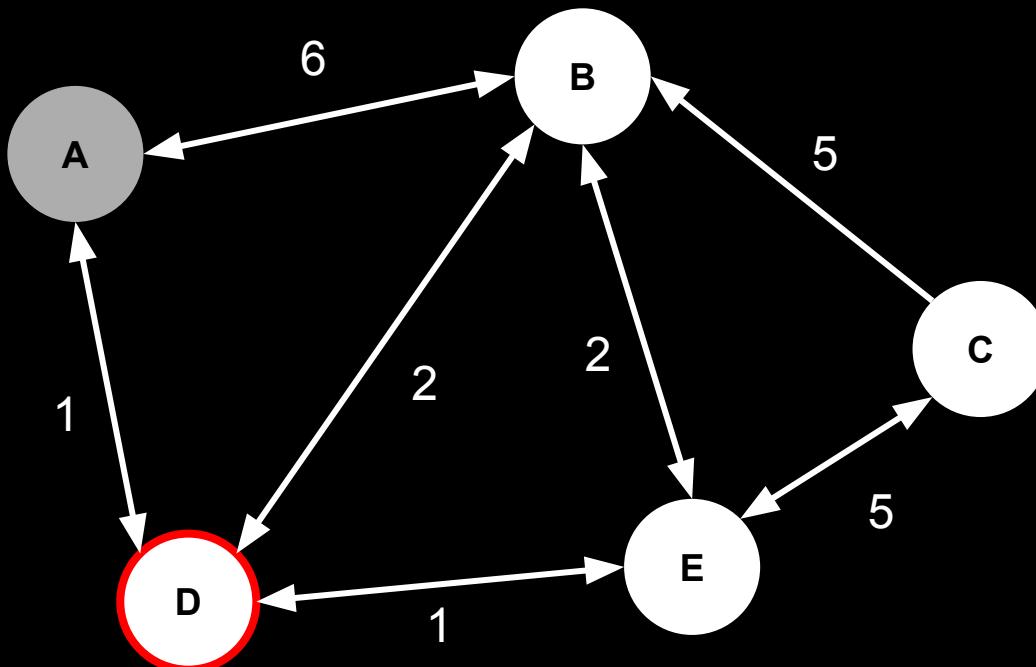
Update distance  
to D to 1



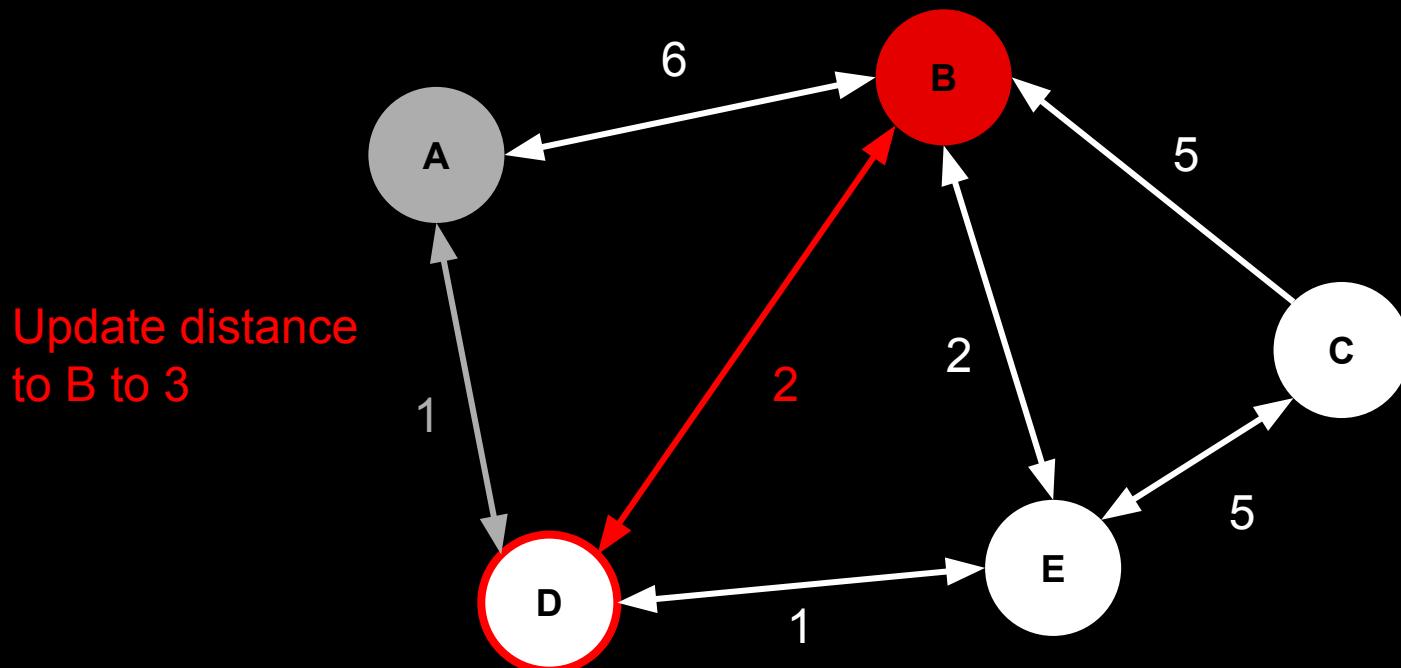
Distances: A = 0, B = 6, C =  $\infty$ , D = 1, E =  $\infty$

Move to D as next location

Mark A as visited

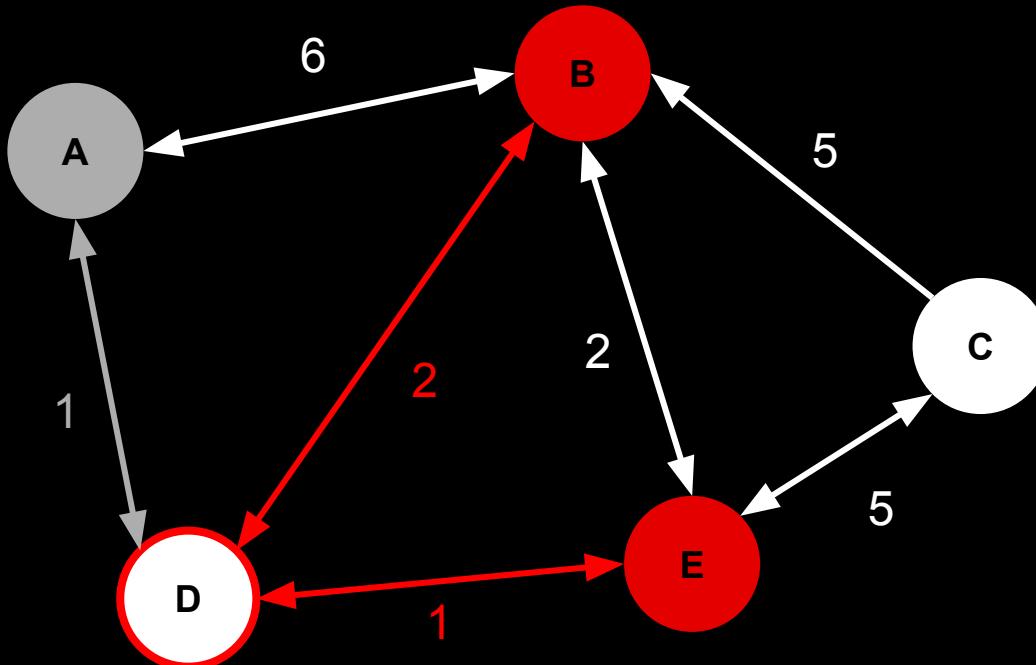


Distances: A = 0, B = 3, C =  $\infty$ , D = 1, E =  $\infty$

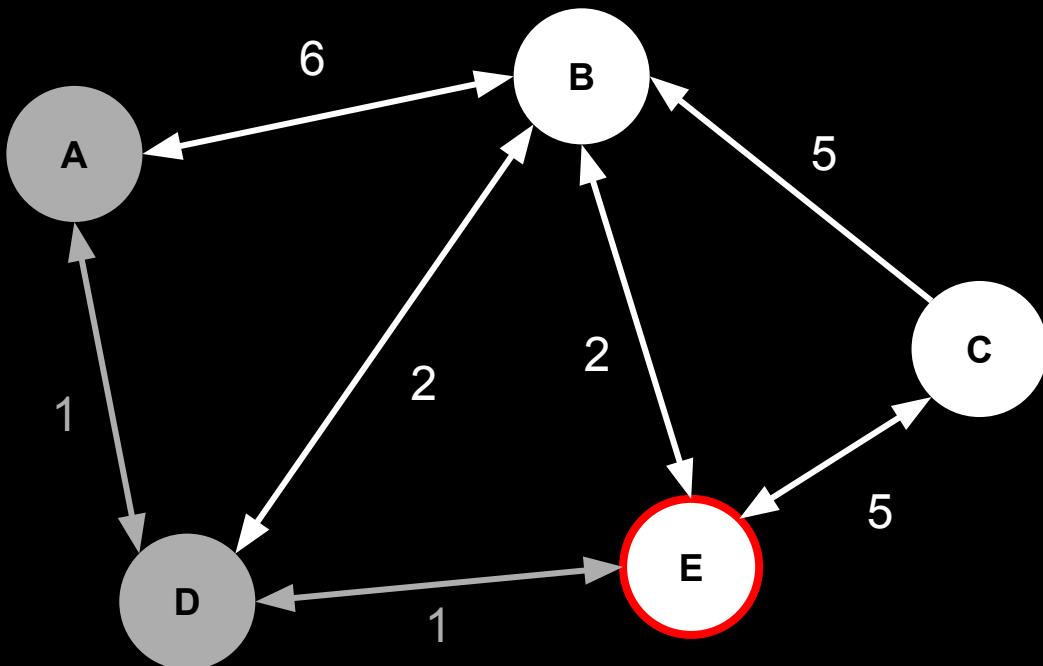


Distances: A = 0, B = 3, C =  $\infty$ , D = 1, E = 2

Update distance  
to E to 2

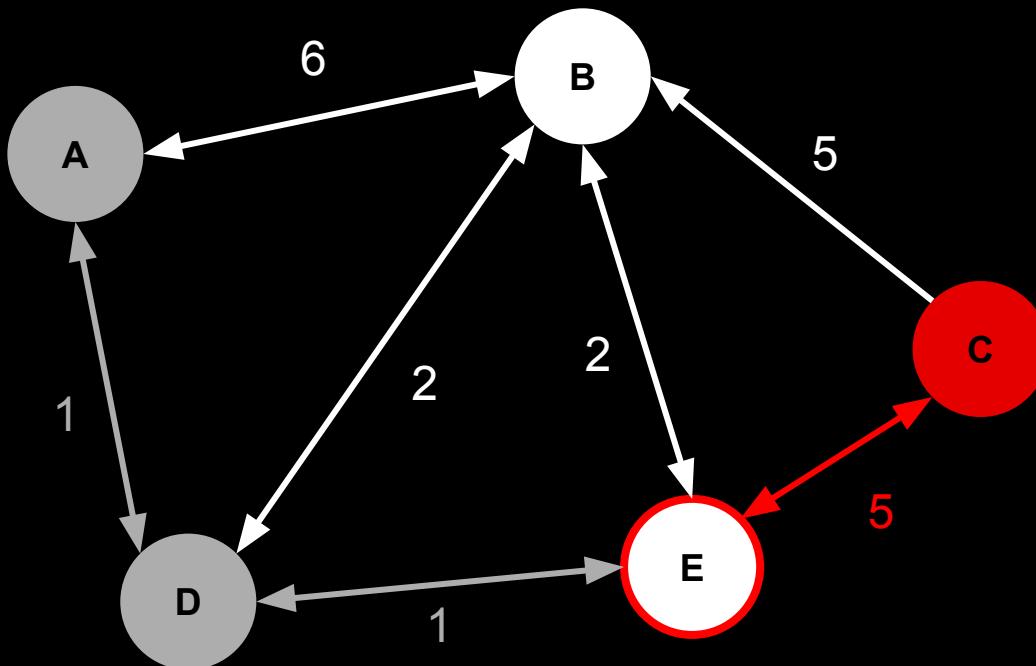


Distances: A = 0, B = 3, C =  $\infty$ , D = 1, E = 2



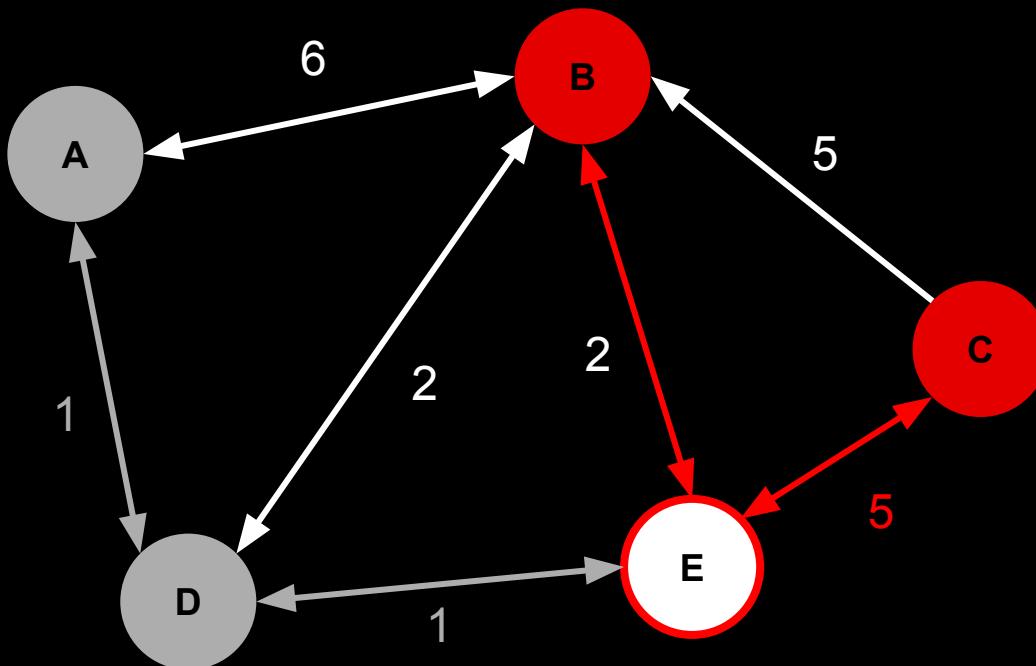
Distances: A = 0, B = 3, C = 7, D = 1, E = 2

Update distance  
to C to 7



Distances: A = 0, B = 3, C = 7, D = 1, E = 2

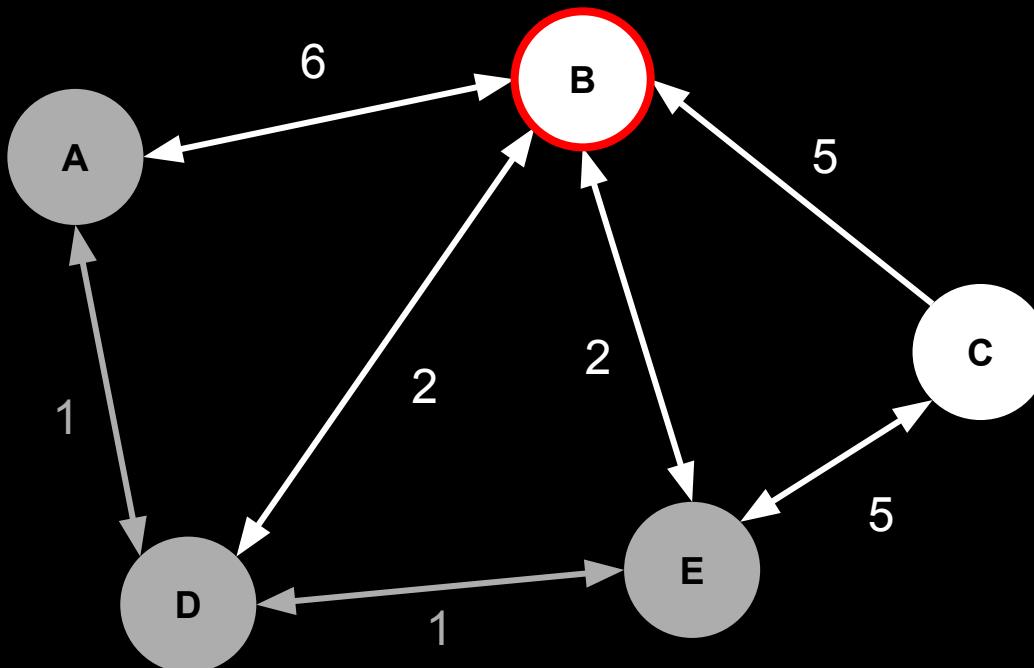
No update for B,  
shorter path  
exists



Distances: A = 0, B = 3, C = 7, D = 1, E = 2

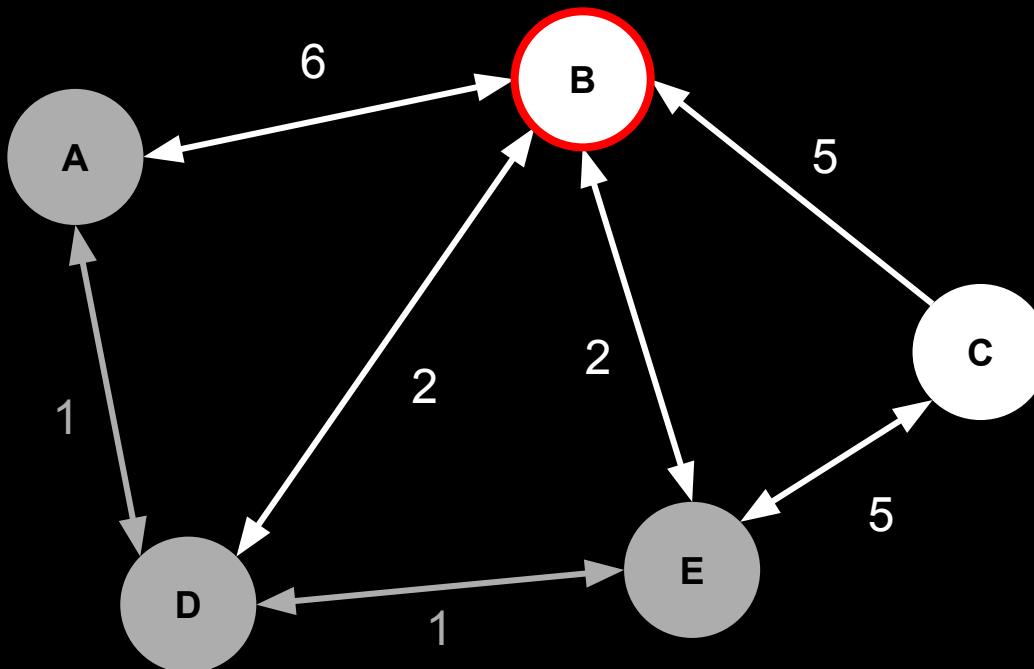
Move to B as  
new location

Mark E as  
visited



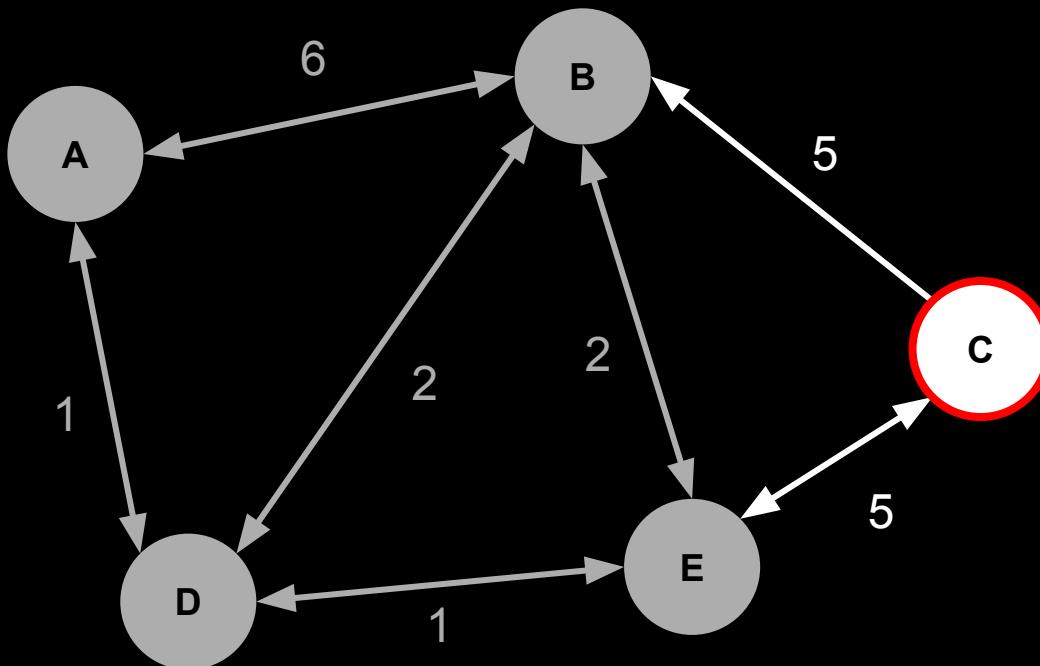
Distances: A = 0, B = 3, C = 7, D = 1, E = 2

No further locations  
reachable  
from B



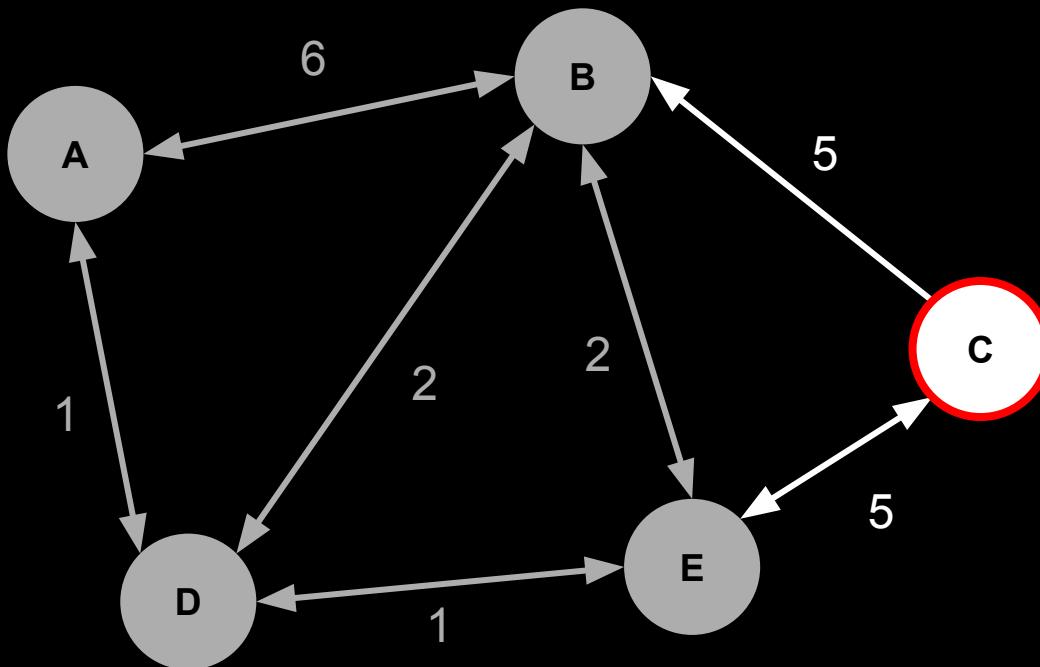
Distances: A = 0, B = 3, C = 7, D = 1, E = 2

Move to C as  
new location



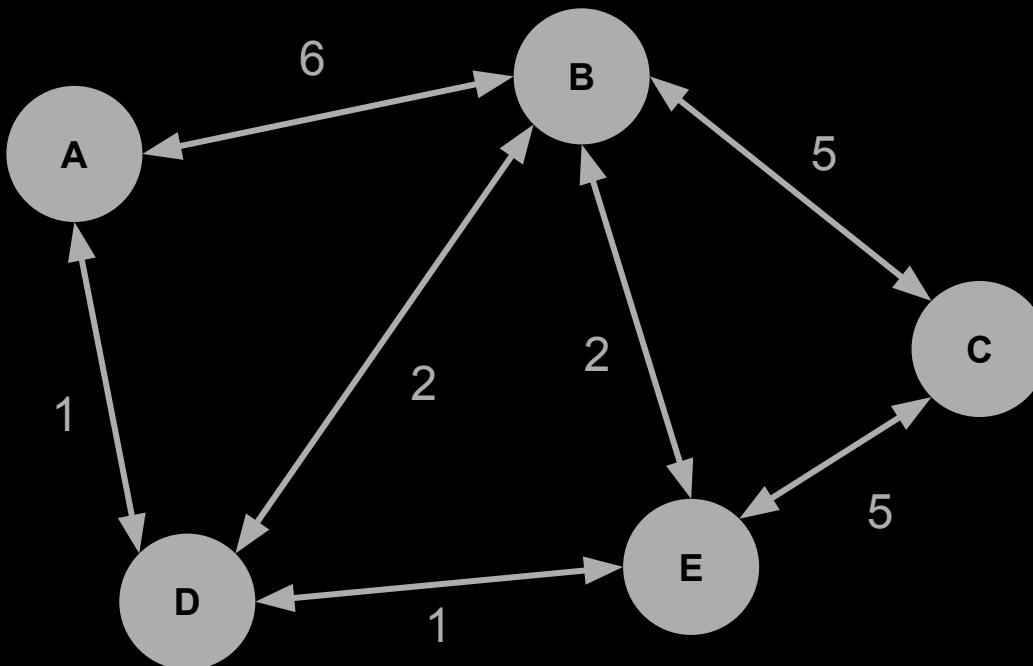
Distances: A = 0, B = 3, C = 7, D = 1, E = 2

No further  
locations  
reachable from  
C



Distances: A = 0, B = 3, C = 7, D = 1, E = 2

All nodes  
visited, we're  
done!





# spam emails



# finding oranges in images

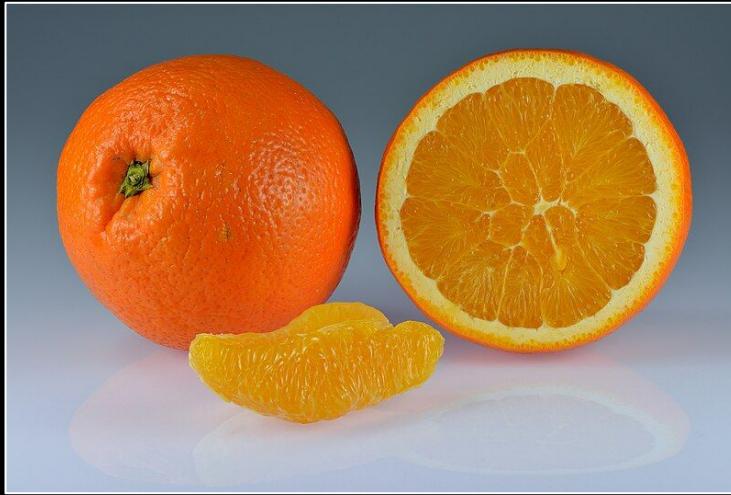


Image source: [Wikimedia](#)

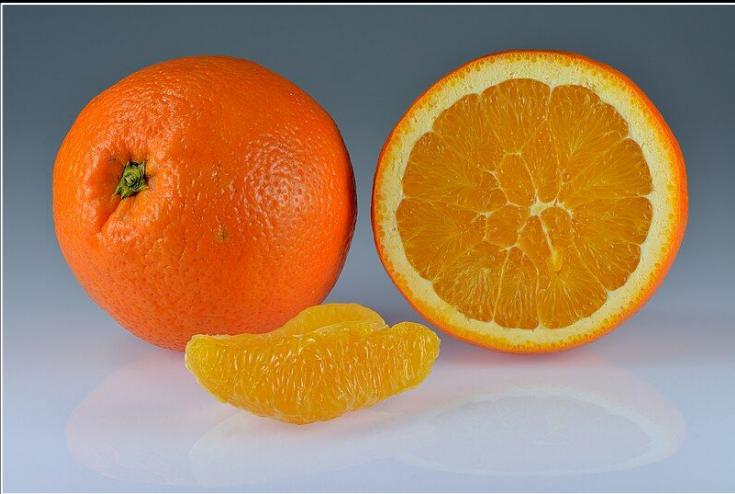


Image source: [Wikimedia](#)

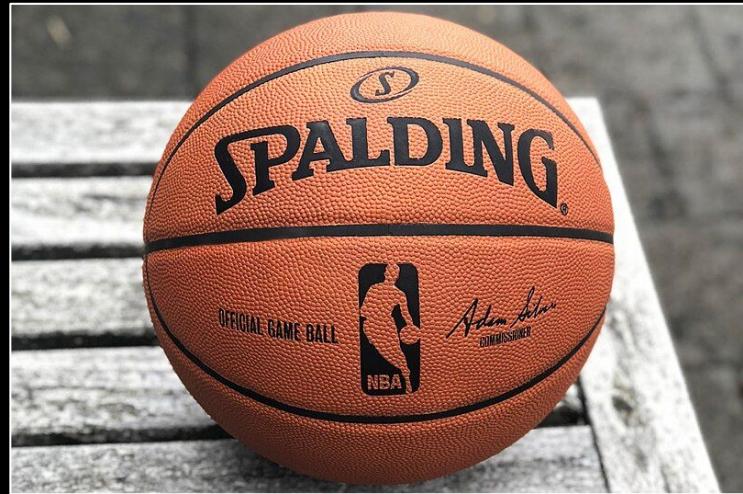


Image source: [Wikimedia](#)

what set of rules can solve this?

machine learning algorithms

