

Lab 3 Writeup

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This lab took me about 10 hours to do. I did attend the lab session.

1. Program Structure and Design:

TCP Sender 功能: 输出的字节流 --> 不可靠数据报的有效负载段

1. 填充, 发包

- * 根据 Sender 当前状态对可发送的窗口

2. 确认对方当前收到的字节流进度

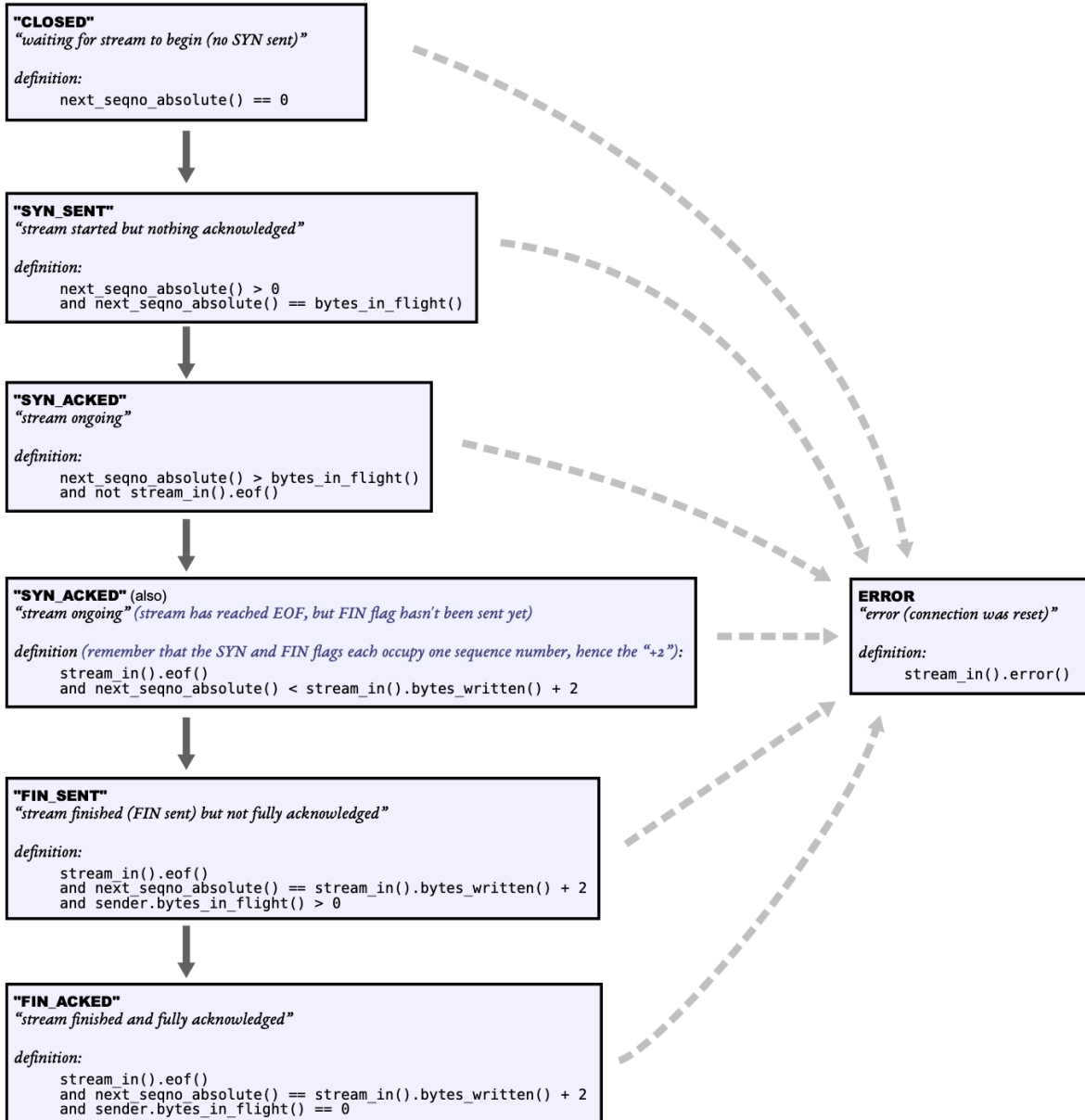
- * 根据对方通知的窗口大小和 ackno

3. 超时重传机制

- * 根据时间变化 (RTO), 定时重传还没有 ack 的报文

<u>3.1 void fill_window() 窗口填充</u>

思路 & 核心代码:



* 配合重传机制的实现：当填充成功一个数据包时，也需要对其进行缓存备份。

(1) //CLOSED

- 条件：初始化
- 需要发送 SYN 包
- 需要填充 SYN 包

Code:

```
//CLOSED
if(next_seqno_absolute() == 0)
{
```

```

//Send the syn
TCPSegment seg;
seg.header().syn = true;
seg.header().seqno = _isn;

_segments_out.push(seg);

_next_seqno = _next_seqno + seg.length_in_sequence_space();
bytes_in_flight_ = bytes_in_flight_ + seg.length_in_sequence_space();

//Cache the segments
_cache_segment(next_seqno_absolute() , seg);
}

```

(2) //SYN_ACKED

- 条件: `ByteStream` 中还有数据可写入 Sender , 且对方窗口大小足够
- 正常按照 payload 大小限制填充数据包

Code:

```

//SYN_ACKED
if(!stream_in().eof() && next_seqno_absolute() > bytes_in_flight())
{
    size_t seg_maxsz_ = TCPConfig::MAX_PAYLOAD_SIZE;

    size_t wdws_remsz_ = peer_windows_size_ ? peer_windows_size_ : 1;
    size_t flight_size_ = bytes_in_flight();
    if(wdws_remsz_ < flight_size_)
        return;

    wdws_remsz_ = wdws_remsz_ - flight_size_;
    string read_stream_ = _stream.read(min(seg_maxsz_, wdws_remsz_));

    while(!read_stream_.empty())
    {
        TCPSegment seg;
        seg.header().seqno = next_seqno();
        seg.payload() = Buffer(std::move(read_stream_));

        wdws_remsz_ = wdws_remsz_ - seg.length_in_sequence_space();

        if(stream_in().eof() && next_seqno_absolute() <
stream_in().bytes_written() + 2 && wdws_remsz_ >= 1)
        {
            seg.header().fin = true;

```

```

        wdws_remsz_ = wdws_remsz_ - 1;
    }

    _next_seqno = _next_seqno + seg.length_in_sequence_space();
    bytes_in_flight_ = bytes_in_flight_ + seg.length_in_sequence_space();
    _segments_out.push(seg);

    //Cache the seg
    _cache_segment(next_seqno_absolute(), seg);
    read_stream_ = _stream.read(min(seg_maxsz_, wdws_remsz_));
}
}

```

(3) //SYN_ACKED

- 条件: `ByteStream`` 已经 ``eof``, 但是 FIN 包还未发送
- 需要填充 FIN 包

Code:

```

//SYN_ACKED
//stream has reached EOF, but FIN flag hasn't been sent yet
if(stream_in().eof() && next_seqno_absolute() < stream_in().bytes_written() +
2)
{
    //Send the fin
    size_t wdws_remsz_ = peer_windows_size_ ? peer_windows_size_ : 1;
    size_t flight_size_ = bytes_in_flight();
    if(wdws_remsz_ <= flight_size_)
        return;

    TCPSegment seg;
    seg.header().seqno = next_seqno();
    seg.header().fin = true;

    _next_seqno = _next_seqno + seg.length_in_sequence_space();
    bytes_in_flight_ = bytes_in_flight_ + seg.length_in_sequence_space();
    windows_size_ = windows_size_ - seg.length_in_sequence_space();

    _segments_out.push(seg);

    //Cache the segments
    _cache_segment(next_seqno_absolute(), seg);
}

```

<u>3.2 void ack_received(const WrappingInt32 ackno, const uint16_t window_size)
ACK 确认</u>

思路 & 核心代码:

- 若为有效确认（缓存的未确认数据包被成功确认），则需要：
 - 更新重传机制相关变量（重传时间，重传次数）
 - 删除确认成功数据包的缓存
 - 填充窗口 -- 因为对端成功确认了数据，对端可用窗口变大了
- 更新对端窗口大小

Code:

```
int n_bytes_ = next_seqno() - ackno;
if(n_bytes_ < 0)
    return;

//Pop all the cache segments
vector<uint64_t> acknos_pop_;

for(auto & seg : segments_cache_)
{
    uint64_t abs_ackno_ = unwrap(ackno, _isn, seg.first);
    if(abs_ackno_ >= seg.first)
        acknos_pop_.push_back(seg.first);
}

for(auto & ackno_pop_ : acknos_pop_)
    segments_cache_.erase(ackno_pop_);

//Extend window size
if(window_size > peer_windows_size_)
{
    windows_size_ += window_size - peer_windows_size_;
}
peer_windows_size_ = window_size;

//Valid ackno
if(!acknos_pop_.empty())
{
    windows_size_ = max(window_size, uint16_t(1));
    bytes_in_flight_ = n_bytes_;

    time_ = 0;
    retx_times_ = 0;
```

```
        _initial_retransmission_timeout =  
default_initial_retransmission_timeout;
```

```
        fill_window();  
    }
```

<u>3.3 void tick(const size_t ms_since_last_tick): 定时重传</u>

思路 & 核心代码:

- 每发送数据包会对其进行缓存
- 当超过 RTO 时间没有收到这个包的确认时,
 - 执行重传
 - 更新等待时间为当前 $RTO * 2$,
- 若有多个数据包超时, 只重传 `seqno` 最小的数据包

Code:

```
    if(segments_cache_.empty())  
        return;  
  
    time_ = time_ + ms_since_last_tick;  
  
    if(time_ >= _initial_retransmission_timeout && retx_times_ <=  
TCPConfig::MAX_RETX_ATTEMPTS)  
    {  
        _segments_out.push(segments_cache_.begin()->second);  
        time_ = 0;  
        if(peer_windows_size_ != 0)  
        {  
            _initial_retransmission_timeout = _initial_retransmission_timeout *  
2;  
            retx_times_ += 1;  
        }  
    }
```

2. Implementation Challenges:

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3. Remaining Bugs:

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