

# Problem Statement :

- Given :
  - A cost matrix `cost[][]` (All costs are positive)
  - Position `(m, n)` in `cost[][]`
- Write a function that returns cost of minimum cost path to reach `(m, n)` from `(0, 0)`.
  - Total cost of a path to reach `(m, n)` is sum of all the costs on that path (including both source and destination)

# Allowed movements :

Can only traverse **up**, **right** and **diagonally upper** cells from a given cell.

❖ From a given cell  $(i, j)$

- $(i+1, j)$
- $(i, j+1)$
- $(i+1, j+1)$

2	1	5	3
1	4	8	2
0	1	2	3
	0	1	2

1	5	3
4	8	2
1	2	3



Minimum cost :

$$1 + 2 + 2 + 3 = 8$$

# Optimal Substructure :

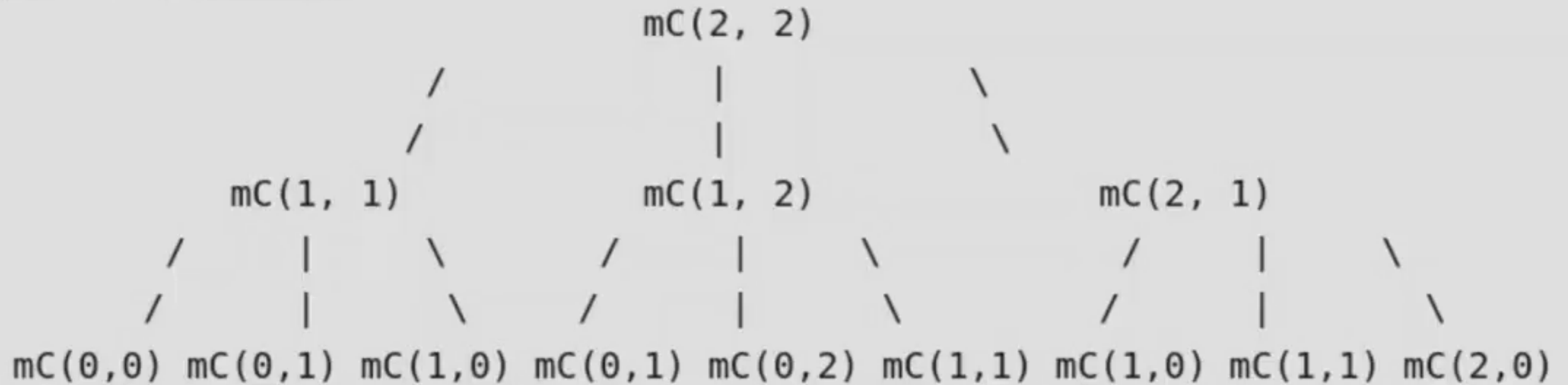
- The path to reach  $(m, n)$  :
  - ◆ Through one of the 3 cells:  $(m-1, n-1)$  or  $(m-1, n)$  or  $(m, n-1)$ .
- $\text{minCost}(m, n) = \min (\text{minCost}(m-1, n-1), \text{minCost}(m-1, n), \text{minCost}(m, n-1))$   
+  $\text{cost}[m][n]$



Time complexity of this naive recursive solution is exponential !

# Overlapping Subproblems

mC refers to minCost()





**So the problem has :**

Optimal Substructure  
and  
Overlapping subproblems

Dynamic Programming

Constructing a temporary array `tc[][]`  
in bottom up manner.

2	1	5	3
1	4	8	2
0	1	2	3
	0	1	2

Cost

2	1	5	3
1	4	8	2
0	1	2	3
	0	1	2

tc

2			
1			
0			
	0	1	2

**i**

**j**

```
int minCost(int cost[R][C], int m, int n)
{
    int i, j;

    // Instead of following line, we can use int tc[m+1][n+1] or
    // dynamically allocate memory to save space. The following line is
    // used to keep the program simple and make it working on all compilers
    int tc[R][C];

    tc[0][0] = cost[0][0];

    /* Initialize first column of total cost(tc) array */
    for (i = 1; i <= m; i++)
        tc[i][0] = tc[i-1][0] + cost[i][0];

    /* Initialize first row of tc array */
    for (j = 1; j <= n; j++)
        tc[0][j] = tc[0][j-1] + cost[0][j];

    /* Construct rest of the tc array */
    for (i = 1; i <= m; i++)
        for (j = 1; j <= n; j++)
            tc[i][j] = min(tc[i-1][j-1],
                           tc[i-1][j],
                           tc[i][j-1]) + cost[i][j];

    return tc[m][n];
}
```

# Time Complexity

$$O(m * n)$$



Thank You for watching !  
Please leave us your likes and comments.

