

# What I learned about functional programming while writing a book on it

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# Why write a book about functional programming. II

I found the FP community to be unlike other programmers' communities

- Others are focused on a chosen programming language (Java, Python, JavaScript, etc.), and on designing and using libraries and frameworks
  - ▶ *“use this framework, override this method, use this annotation”*
- The FP community talks in a very different way
  - ▶ *“referential transparency, algebraic data types, monoid laws, parametric polymorphism, free applicative functors, monad transformers, Yoneda lemma, Curry-Howard isomorphism, profunctor lenses, catamorphisms”*
    - ★ A glossary of FP terminology (more than 100 terms)
  - ▶ From SBTB 2018: *The Functor, Applicative, Monad talk*
    - ★ By 2018, everyone expects to hear a talk about these concepts
  - ▶ An actual Scala error message:

```
found    : Seq[Some[V]]  
required: Seq[Option[?V8]] where type ?V8 <: V (this is a GADT skolem)
```

To do FP, should I learn all of this? How do I learn about this?

# Why write a book about functional programming. III

Main questions:

- Which theoretical knowledge will actually help write Scala code?
- Where can one learn this FP theory, with definitions and examples?

What I did *not* want to see:

- Heuristic explanations without derivations and proofs
  - ▶ Most FP books show code without proofs or rigorous definitions
    - ★ *The Book of Monads* does not prove the laws for any monads
  - ▶ A few books (*Haskell Wikibooks*, *Introduction to functional programming using Haskell*, and *Functional programming in Scala*) include some simple proofs
- Academic theory for the sake of theory, with no applications
  - ▶ “*Monad is just a monoid in the category of endofunctors*”
  - ▶ *The Book of Monads* chapter 18 (adjunctions)
  - ▶ *Lawvere theories* as an alternative to monads

# Why write a book about functional programming. IV

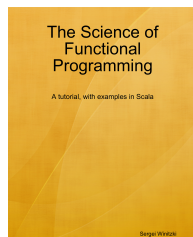
Reading various materials has given me more questions than answers  
I started writing a **new book** to answer all my FP questions

- by motivating and deriving all results from scratch
- organizing systematically the practice-relevant parts of FP theory

The book explains (with code examples and exercises):

- theory and applications of major design patterns of FP
- techniques for deriving and verifying properties of types and code (typeclass laws, equivalence of types)
- practical motivations for (and applications of) these techniques

Status of the book: 12.5 out of 14 chapters are ready



# What I learned. I. Questions that have rigorous answers

In FP, a programmer encounters certain questions about code that can be answered rigorously

- The answers will guide the programmer in designing the code
- The answers are *not* a matter of opinion or experience
- The answers are found via mathematical derivations and reasoning

# Examples of reasoning tasks. I

- 1 Can we compute a value of type `Either[Z, R => A]` given a value of type `R => Either[Z, A]` and conversely? (`A`, `R`, `Z` are type parameters.)

```
def f[Z, R, A](r: R => Either[Z, A]): Either[Z, R => A] = ???  
def g[Z, R, A](e: Either[Z, R => A]): R => Either[Z, A] = ???
```

- We can implement `g`, and there is only one way:

```
def g[Z, R, A](e: Either[Z, R => A]): R => Either[Z, A] =  
  r => e.map(f => f(r))           // Scala 2.12 code
```

- It turns out that `f` *cannot* be implemented
  - ▶ Not because we are insufficiently clever, but because... math!
- Programmers need to develop intuition about why this is so
- These results are rigorous
  - ▶ The Curry-Howard isomorphism and the LJT algorithm
  - ▶ The code for `g[Z, R, A]` can be generated automatically

# Examples of reasoning tasks. II

- 2 How to use `for / yield` with `Either[Z, A]` and `Future[A]` together?

```
val result = for { // This code will not compile; need to combine...  
  a <- Future(...) // ... a computation that is run asynchronously,  
  b <- Either(...) // a computation whose result may be unavailable,  
  c <- Future(...) // and another asynchronous computation.  
} yield ???      // Continue computations when results are available.
```

Should `result` have type `Either[Z, Future[A]]` or `Future[Either[Z, A]]`?

How to combine `Either` with `Future` so that we can use `flatMap`?

- It turns out that `Either[Z, Future[A]]` is wrong (cannot implement `flatMap` correctly). The correct type is `Future[Either[Z, A]]`.
- Programmers need to develop intuition about why this is so
- This is a rigorous result (programmers do not need to test it)
  - ▶ The theory of monad transformers and their laws



# Examples of reasoning tasks. III

- ③ Can we implement `flatMap` for the type constructor `Option[(A, A, A)]`?

```
def flatMap[A, B](fa: Option[(A, A, A)])(f: A => Option[(B, B, B)])  
  : Option[(B, B, B)] = ???
```

- It turns out that `flatMap` *can* be implemented but fails the monad laws
- Programmers need to develop intuition about why this is so
  - ▶ How should we modify `Option[(A, A, A)]` to make it into a monad?
- This is a rigorous result (programmers do not need to test it)
  - ▶ The theory of monads and their laws
  - ▶ The theory of type constructions of monads

# Examples of reasoning tasks. IV

- ④ Different people define a “free monad” via different sets of case classes. Are these definitions equivalent? What is the difference?
  - ▶ Three different implementations of the free monad: a **blog post** by G. Gonzalez (2012), a **talk** given by R. Bjarnason (2014), and a **talk** given by K. Robinson (2016) — but no rigorous definitions
- The free monad on a functor is less code than the free monad on a non-functor
- The free monad's encoding that assumes the monad laws is less code than an encoding without assumed laws
- These are rigorous results (programmers do not need to test them)
  - ▶ The theory of “free” inductive typeclasses and their encodings
- Programmers need intuition about implementing free typeclasses
  - ▶ Given a **Pointed** functor (a functor that already has the **pure** method), implement a free monad or a free applicative

# What I learned. II. Functional programming is engineering

- FP is similar to engineering in some ways
  - ▶ Mechanical, electrical, chemical engineering are based on calculus, classical and quantum mechanics, electrodynamics, thermodynamics
    - ★ These sciences give engineers rigorous answers to certain questions relevant to engineering design
  - ▶ FP is based on category theory, type theory, logic proof theory
    - ★ These theories give programmers rigorous answers to certain questions relevant to writing code
  - ▶ Programming in non-FP paradigms is similar to *artisanship*
- Engineers use special terminology
  - ▶ Examples from mechanical, electrical, chemical engineering: rank-4 tensors, Lagrangians with non-holonomic constraints, Fourier transform of the delta function, inverse Z-transform, Gibbs free energy
  - ▶ Examples from FP: rank- $N$  types, continuation-passing transformation, polymorphic lambda functions, free monads, hylomorphisms
- As in engineering, the special terminology in FP is *not* self-explanatory
  - ▶ What is a delta function? What is a lambda function?
  - ▶ What is the Gibbs free energy? What is the free monad?

# What I learned. III. The science of map / filter / reduce

The `map/filter/reduce` (MFR) programming style: iteration without loops

- Compute the list of all integers  $n$  between 1 and 100 that can be expressed as  $n = p * q$  (with  $2 \leq p \leq q$ ) in exactly 4 different ways

```
scala> (1 to 100).filter { n =>
  |   4 == (2 to n).count { x => n % x == 0 && x * x <= n }
  | }
res0: IndexedSeq[Int] = Vector(36, 48, 80, 100)
```

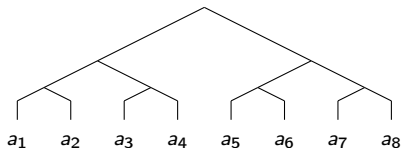
The MFR programming style is an FP success story

- Nameless functions (“lambdas”, “closures”) are widely used
  - ▶ and have been added to most programming languages by now
- Essential methods: `map`, `filter`, `flatMap`, `zip`, `fold`
- Similar techniques work with parallel and stream processing (`Spark`)
- Similar techniques work with relational databases (`Slick`)

# What I learned. III. The science of map / filter / reduce

Essential MFR methods: `map`, `filter`, `flatMap`, `zip`, `fold`

- What data types other than `Seq[A]` can support these methods?
  - ▶ Algebraic data types
  - ▶ Trees and other recursive types
  - ▶ Perfect-shaped trees



- ▶ Which methods can be defined for `MyData[A]`?  
`type MyData[A] = String => Option[(String, A)]`

# What I learned. III. The science of map / filter / reduce

A systematic approach to understanding FP via a study of MFR

- Determine the required laws of `map`, `filter`, `flatMap`, `zip`, `fold`
  - ▶ The laws express the programmers' expectations about code behavior
  - ▶ Define the corresponding typeclasses
    - ★ `map` — `Functor`, `filter` — `Filterable`, `flatMap` — `Monad`, `zip` — `Applicative`, `fold` — `Traversable`
- Find type constructions that preserve the typeclass laws
  - ▶ If `P[A]` and `Q[A]` are filterable functors then so is `Either[P[A], Q[A]]`
  - ▶ If `P[A]` is a monad then so is `Either[A, P[A]]`
  - ▶ If `P[A]` and `Q[A]` are monads then so is `(P[A], Q[A])`
  - ▶ If `P[A]` is a contravariant functor then `P[A] => A` is a monad
  - ▶ If `P[A]` and `Q[A]` are applicative then so is `Either[P[A], (A, Q[A])]`
    - ★ I found many more type constructions of this kind
  - ▶ Sometimes it becomes necessary to define additional typeclasses
    - ★ Contravariant functor, contravariant filterable, contravariant applicative
- Develop intuition about implementing lawful typeclass methods
- Develop intuition about data types that can have those methods
  - ▶ ... and about data types that *cannot* (and reasons why)
- Develop notation and proof techniques for proving the laws

# What I learned. IV. The logic of types

FP is not just “programming with functions”: types play a central role

- The compiler needs to check all types at compile time
- The language needs to support certain type constructions

Most of FP use cases are based on only six type constructions:

- Unit type — `Unit`
- Type parameters — `f[A](x)`
- Product types — `(A, B)`
- Co-product types (“disjunctive union” types) — `Either[A, B]`
- Function types — `A => B`
- Recursive types — `Fix[A, S]` where `S[_ , _]` is a “recursion scheme”  
`final case class Fix[A, S[_ , _]](unfix: S[A, Fix[A, S]])`

Going through all possible type combinations, we can enumerate essentially all possible typeclass instances

- all possible functors, filterables, monads, applicatives, traversables, etc.
- in some cases, we can generate typeclass instances automatically

# What I learned. IV. The logic of types

- Unit, product, co-product, and function types correspond to logical propositions (`true`),  $(A \text{ and } B)$ ,  $(A \text{ or } B)$ ,  $(\text{if } A \text{ then } B)$
- Not all programming languages support all of these type constructions
  - ▶ The logic of types is *incomplete* in those languages
- Languages that do not support co-products will make you suffer

```
fileOpened, err := os.Open("filename.txt")    // go-lang has you  
if err != nil { log.Fatal(err) } // doomed to write this forever
```

- Returning a pair (both a result and an error) instead of a disjunction (either a result or an error) promotes many ways of making hard-to-find mistakes
  - ▶ In Scala, we may just return `Try[Result]` or `Either[Error, Result]`



## What I learned. V. Miscellaneous surprises

My approach forced me to formulate and prove every statement  
Each chapter gave me at least one surprise

- What I believed and tried to prove turned out to be incorrect
- What seemed to be intuitively unexpected turned out to be true

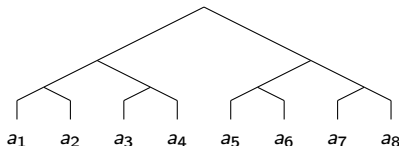
# What I learned. V. Miscellaneous surprises

## Chapters 1 to 3:

- Nameless functions are used in mathematics too, just hidden

$\sum_{n=1}^{100} n^2$	<code>(1 to 100).map { n =&gt; n * n }.sum</code>
$\int_0^1 \sin(x^3) dx$	<code>integrateNumerically(0, 1) { x =&gt; math.sin(x * x * x) }</code>

- Many algorithms require non-tail-recursive code (`map` for a tree)
- Perfect-shaped trees *can* be defined via recursive ADTs



# What I learned. V. Miscellaneous surprises

Chapters 4 and 5: a practical application of the Curry-Howard isomorphism

- “Type inference” — determining type signature from given code
- “Code inference” — determining code from given type signature
- The `curryhoward` library uses the LJT algorithm for code inference

```
import io.chymyst.ch._
```

```
scala> def in[A, B](a: A, b: Option[B]): Option[(A, B)] = implement
def in[A, B](a: A, b: Option[B]): Option[(A, B)]
```

```
scala> in(1.5, Some(true))
val res0: Option[(Double, Boolean)] = Some((1.5,true))
```

```
scala> def h[A, B]: (((A => B) => A) => A) => B = implement
def h[A, B]: (((A => B) => A) => A) => B
```

```
scala> println(h.lambdaTerm.prettyPrint)
a => a (b => b (c => a (d => c)))
```

```
scala> def g[A, B]: (((A => B) => B) => A) => B = implement
error: type (((A => B) => B) => A) => B cannot be implemented
```

# What I learned. V. Miscellaneous surprises

## Chapters 6 to 8:

- Functions of type  $\text{ADT} \Rightarrow \text{ADT}$  can be manipulated via matrices
  - ▶ Matrix code notation is useful in symbolic proofs

```
val p: Either[A, B] => Either[C, D] = {  
  case Left(x)    => Right(f(x))  
  case Right(y)   => Left(g(y))  
}
```

	C	D
A	0	$x \rightarrow f(x)$
B	$y \rightarrow g(y)$	0

- Typeclasses can be viewed as partial functions from types to values
- All non-parameterized types have a monoid structure
- Subtypes / supertypes are not always the same as supersets / subsets

# What I learned. V. Miscellaneous surprises

## Chapters 9 to 12:

- “Filterable functors” are a neglected typeclass with useful properties
- Data types `Option[(A, A)]`, `Option[(A, A, A)]`, etc., *cannot* be monads
- Monads need “runners” to be useful, but some monads’ runners do not obey the laws or cannot exist (`State`, `Continuation`)
- Without some laws, `flatMap` is *not* equivalent to `map` with `flatten`
  - ▶ It is not enough to write `_.flatten == _.flatMap(identity)` and `_.flatMap(f) == _.map(f).flatten`, we need to prove an isomorphism
- All contravariant functors are applicative (if defined using the six standard type constructions)
- Breadth-first traversal of trees *can* be defined via `fold` and `traverse` (not only depth-first traversal)

# What I learned. V. Miscellaneous surprises

## Chapter 13 (free typeclass constructions):

- Not all typeclasses have a “free” construction: there is free functor, filterable, applicative, etc.; but *no* free foldable or free traversable
- *“Tagless final” is just a Church encoding of the free monad, what is the problem?*
- It is hard to prove the correctness of the Church encoding
  - ▶ My book uses relational parametricity together with some results from **unpublished talk slides** to prove that the Church encoding works
  - ▶ ... but programmers do not need to study those proofs

## Chapter 14 (monad transformers):

- Monad transformers likely exist for all explicitly definable monads, but there is *no* general method or scheme for defining the transformers
- Some monad transformers are incomplete, not fully usable for combining monadic effects ([Continuation](#), [Codensity](#))
- *Monad transformers are just pointed endofunctors in the category of monads, what is the problem?*
- Monad transformers have 18 laws

# Conclusions

- Functional programming has a steep learning curve
  - ▶ Programmers can already benefit from the simplest techniques
    - ★ ... and mostly stop there (`map` / `filter` / `fold`, ADTs, `for` / `yield`)
  - ▶ Full *ab initio* derivations and proofs take 500 pages
  - ▶ The difficulty is at the level of undergraduate calculus / algebra
- Much of the theory is directly beneficial for coding
- Using FP techniques makes programmers' work closer to *engineering*
- Full details in the free book — <https://github.com/winitzki/sofp>