

Numerical Analysis Qualifying Exam

Fall 1997

1. Explain the difficulty in evaluating $f(x) = x^{-1} (1 - \cos x)$ when the absolute value of x is small and find a method without using the Taylor series to overcome the difficulty.
2. Prove the following theorem: If $f \in C^2(a, b)$, $f'(x) f''(x) \neq 0$, and $f(x)$ has a zero in (a, b) , then the zero is unique in (a, b) , and the Newton iteration will converge to it if the starting value x_0 and the first approximation x_1 are both in (a, b) . (You may just do a special case where $f'(x) < 0$, $f''(x) < 0$ in (a, b))

3. Prove the following theorem for Gaussian quadrature:

Let $I(f) = \int_a^b f(x)w(x)dx$, where $w(x)$ is a positive weight function, be approximated by a quadrature formula $I_n(f) = \sum_{i=1}^n A_i f(x_i)$, where $x_i \in [a, b]$, $i = 1 : n$ are distinct. Then the quadrature formula $I_n(f)$ has a maximum degree of precision of $2n - 1$. This is attained if and only if x_1, x_2, \dots, x_n are the zeros of $p_n(x)$, the n th orthogonal polynomial, with the inner product

$$(f, g) = \int_a^b f(x)g(x)w(x)dx.$$

(Hint: if $f(x)$ is a polynomial of degree m with $m \geq n$, then $f[x_1, x_2, \dots, x_n, x]$ is a polynomial of degree $m - n$)

4. Let $f(x)$ be a piecewise constant function on $[a, b]$:

$$f(x) = \begin{cases} \alpha & a \leq x \leq c \\ \beta & c < x \leq b \end{cases}$$

where $a < c < b$ and $\alpha \neq \beta$. Let $x_0 < x_1 < \dots < x_n$ be $n + 1$ points and j be an integer with $0 \leq j \leq n - 1$ such that $a < x_0 < \dots < x_j < c < x_{j+1} < \dots < x_n < b$ and that $P_n(x)$ is the polynomial of degree less than or equal to n interpolating f at the $n + 1$ points x_i , $i = 0, \dots, n$. Show that $P_n(x)$ is monotone in the interval $[x_j, x_{j+1}]$. (Hint: Count the number of zeros of $\frac{dP_n(x)}{dx}$.)

5. Suppose $A \in R^{n \times n}$, $b \in R^n$, and $PA = LU$, where P is a permutation matrix, L is lower triangular, and U is upper triangular. A, b, P, L, U are known. Determine the purpose of the following algorithm:

For $j = 1 : k$

solve for y in $Ly = Pb$ overwrite b with y

solve for x in $Ux = b$ overwrite b with x

end

(Hint: consider $k = 1$ first)

6. Let A be an invertible $n \times n$ matrix, and let u and v be two vectors in R^n . Find the necessary and sufficient conditions on u and v in order that the matrix

$$B = \begin{bmatrix} A & u \\ v^t & 0 \end{bmatrix}.$$

be invertible, and give the inverse in terms of A, u, v when it exists. (Hint: multiply B by a suitable matrix to have a simpler matrix to handle)

7. Suppose that an n by n nonsingular matrix A has distinct eigenvalues $\lambda_1, \dots, \lambda_n$ with corresponding right eigenvectors u_1, \dots, u_n and left eigenvectors v_1, \dots, v_n . Suppose

$$c_i = \|u_i\|_2 \|v_i\|_2 / |v_i^T u_i|, i = 1, \dots, n.$$

Show that the solution of $Ax = b$ satisfies

$$\|x\|_2 \leq \|b\|_2 \sum_{i=1}^n \frac{c_i}{|\lambda_i|}.$$

(Hint: write x in terms of u_i)

8. State Schur's theorem and use it to show: Let A be an $n \times n$ complex matrix, it can be similar to a upper triangular matrix with $|\text{off-diag. entries}| \leq \varepsilon$, (Hint: on the result of Schur's, construct suitable diagonal matrices)