

Algorithms for finding all maximal bicliques of a graph

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June 27, 2019

1 Graphs and bicliques

- Basic concept
- Bicliques

2 Algorithms

- Lexicographic Clique Generation algorithm (LEX)
- Maximal biclique enumeration algorithm (MBEA)
- Modular input consensus algorithm (MICA)

3 Example

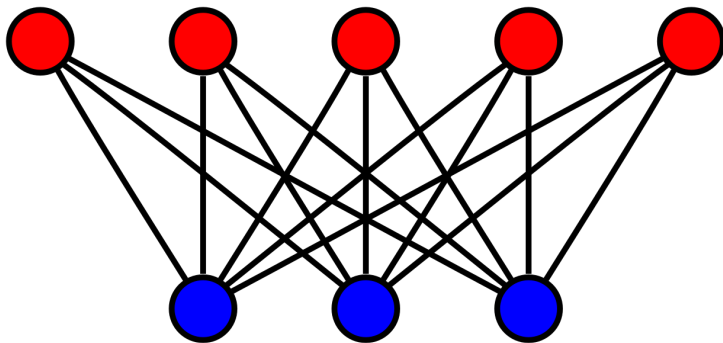
4 Graphical user interface (GUI)

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Basic Concept

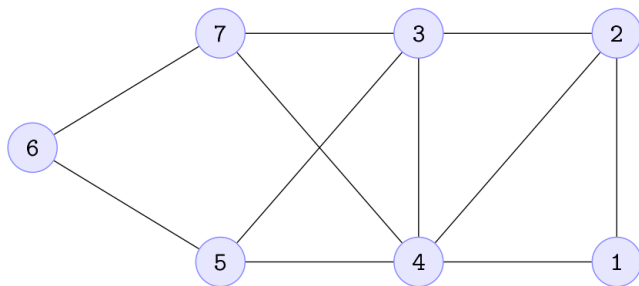
So, there are graphs and they have bicliques...



- Maximum Biclique Problem
- Has many applications, including clustering and m

Example

Compute all maximal bicliques of the following graph:



Algorithm LEX

- The most basic algorithm
- Relies on two biclique operations:
 - Absorption
 - Consensus Adjunction
- Main idea:
 - a seed list of bicliques is created
 - absorption applied for each biclique against every other biclique
 - consensus adjunction applied to create new bicliques
- non polynomial complexity in *both* graph size *and* the number of maximum bicliques!

Algorithm LEX

Let $B_1 = (X_1, Y_1)$ and $B_2 = (X_2, Y_2)$ be two bicliques of G . We say that B_1 *absorbs* or *contains* B_2 if $X_2 \subseteq X_1$ and $Y_2 \subseteq Y_1$, or if $X_2 \subseteq Y_1$ and $Y_2 \subseteq X_1$.

If $Y_1 \cap Y_2 \neq \emptyset$, we call $(X_1 \cup X_2, Y_1 \cap Y_2)$ one of the *consensuses* of B_1 and B_2 . Similarly, each of those pairs of subsets $(X_1 \cap X_2, Y_1 \cup Y_2)$, $(Y_1 \cup X_2, X_1 \cap Y_2)$, $(X_1 \cup Y_2, Y_1 \cap X_2)$ which define bicliques (i.e., which involve two non-empty subsets) are *consensuses* of B_1 and B_2 . In this way, a pair of bicliques may have 0, 1, 2, 3 or 4 consensuses.

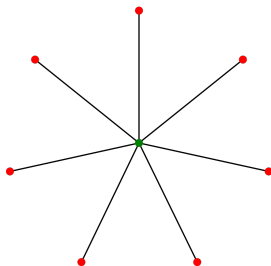
Algorithm MBEA

(not implemented yet)

Algorithm MICA

Modular input consensus algorithm (MICA)

- **Modular consensus:** sequence of operations performed on each pair of bicliques in the repeat loop.
- **Input:** starting with a special input set of maximal bicliques that covers the stars of the graph.



Algorithm MICA

The algorithm has two steps (n is the number of vertices):

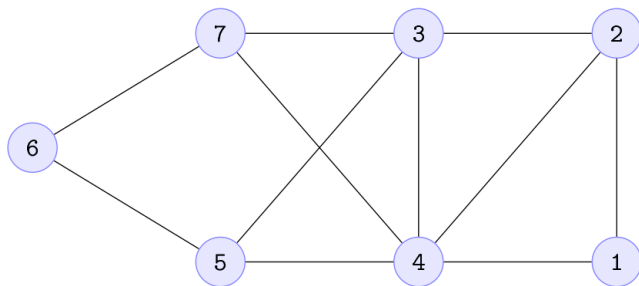
- **Start** with a list C_0 of at most n maximal bicliques that covers the stars of the graph. Let $C := C_0$.
- **Repeat**: For every pair of distinct bicliques $B_1 \in C_0$ and $B_2 \in C$, if B_1 and B_2 have a consensus B_3 which is not absorbed by any member of C , then extend B_3 to a maximal biclique B_4 , and add B_4 to C .

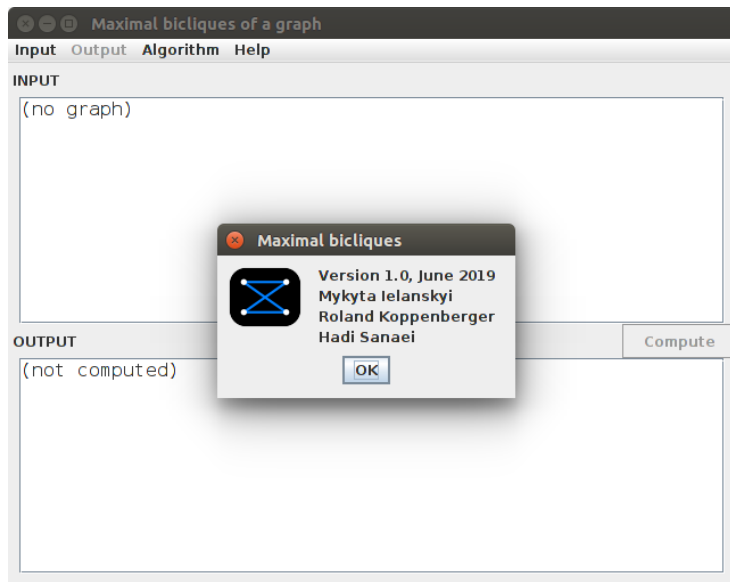
Some substantial improvements for implementation:

- Many consensus adjunction operations produce the same output.
- $X \subset Vertices$: $\Gamma(X) :=$ set of vertices adjacent to each vertex in X .
For all $X \subset Vertices$: $(\Gamma^2(X), \Gamma(X))$ is a maximal biclique.
- Proceed in stages, consider only newly added bicliques in each stage.
- Time complexity: incrementally polynomial, i.e. time for computation of next new maximal biclique is bounded polynomially in number of vertices and number of maximal bicliques computed so far.

Example

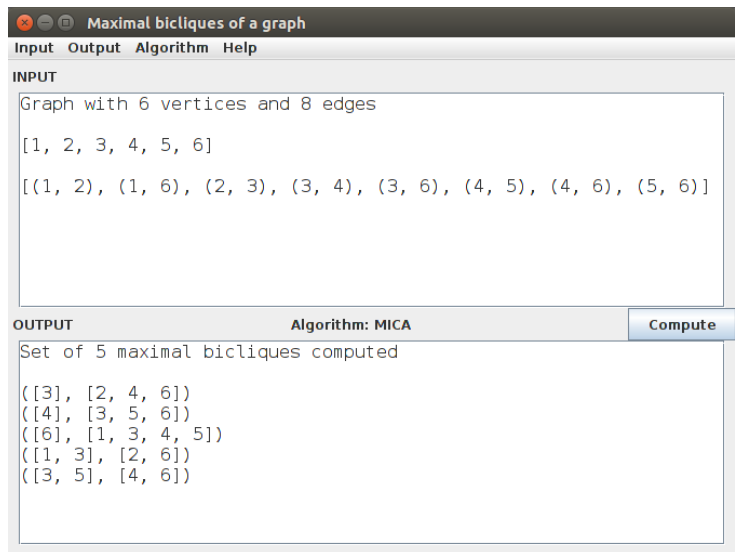
Again: *Compute all maximal bicliques of the following graph:*



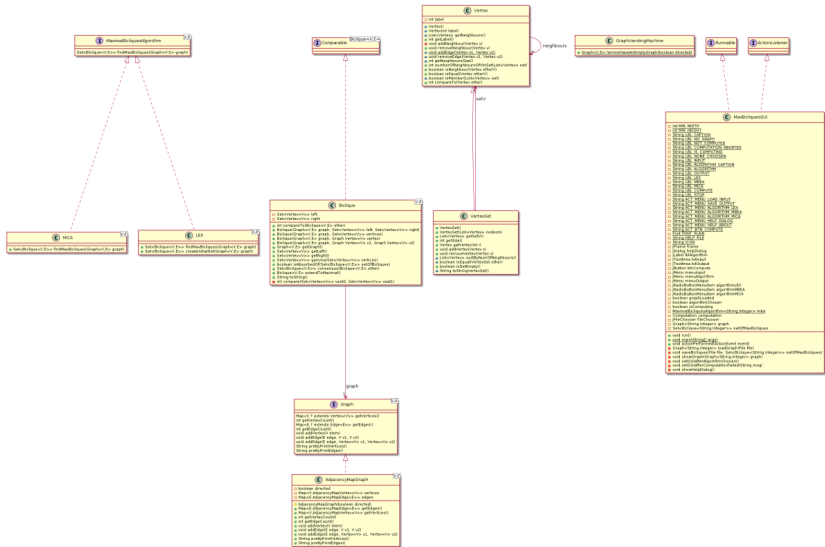


What can one do with the graphical user interface?

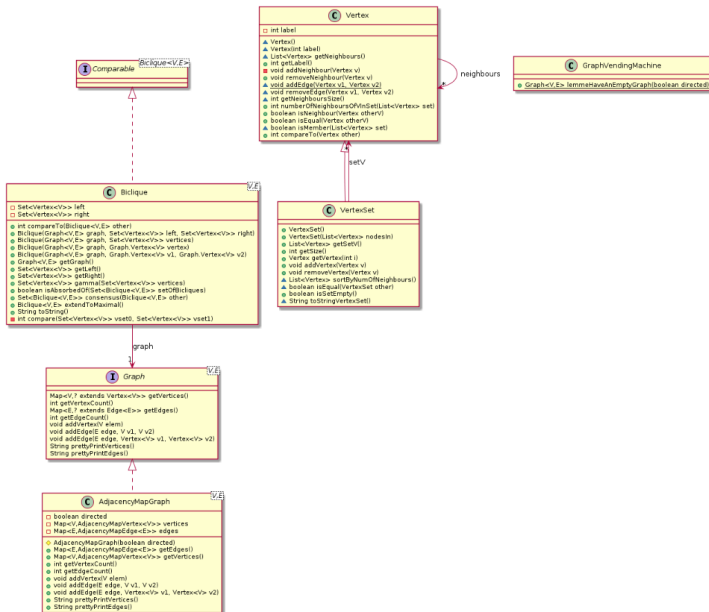
- Load a graph from a text file.
- Select an algorithm for computation.
- Compute the set of maximal bicliques.
- Save the result of the computation in a text file.

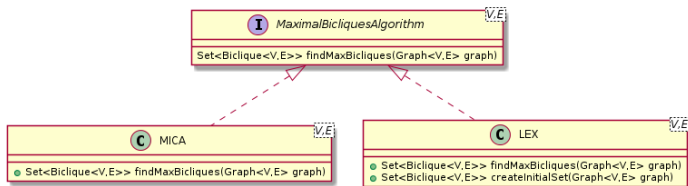


package bicliques



package bicliques.graphs







Practical demonstration

We conclude with some practical computations.