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To classify the languages $L(w)$ given a DFA that accepts a certain language L , we need to check whether the language defined by the tail w (i.e., $(0+1)^*w$) falls into one of three categories:

1. **Contained in L** : All strings ending with w are in L .
2. **Disjoint from L** : No string ending with w is in L .
3. **Neither contained nor disjoint**: Some strings ending with w are in L and some are not.

To make this determination, we follow these steps:

1. **Simulate the DFA**: Check how the DFA processes the string w .
2. **Check the state after processing w** : Determine if the DFA reaches an accepting state after processing w .
3. **Test with arbitrary prefixes**: Check how the DFA reacts to arbitrary prefixes followed by w .

Classification of the Given Languages

1. **$L(1111001)$** :

- Check if the DFA reaches an accepting state after processing the string "1111001".
- If the DFA ends in an accepting state after "1111001", then $L(1111001)$ is contained in L .
- If it does not reach an accepting state, further testing with different prefixes is required.

2. **$L(11011)$** :

- Simulate the DFA on "11011".
- If the DFA ends in an accepting state, then $L(11011)$ is contained in L .
- Otherwise, check with different prefixes to see if any combination of prefixes followed by "11011" is accepted.

3. **$\setminus(L(110101))$** :

- Simulate the DFA on "110101".
- If the DFA ends in an accepting state, then $\setminus(L(110101))$ is contained in $\setminus(L)$.
- If the DFA does not accept "110101", check with various prefixes to determine the relationship with $\setminus(L)$.

4. **$\setminus(L(00011101))$** :

- Simulate the DFA on "00011101".
- If the DFA ends in an accepting state, then $\setminus(L(00011101))$ is contained in $\setminus(L)$.
- If not, test with different prefixes to see if there are any strings in $\setminus(L)$ that end with "00011101".

General Approach

1. **Simulate the DFA**: Process the string $\setminus(w)$ through the DFA.

2. **Analyze the Result**:

- If the DFA reaches an accepting state, $\setminus(L(w))$ is contained in $\setminus(L)$.
- If the DFA does not reach an accepting state, further check is required with various prefixes.

3. **Test Prefixes**:

- If some prefixes followed by $\setminus(w)$ are accepted and some are not, then $\setminus(L(w))$ is neither contained nor disjoint.
- If no prefixes followed by $\setminus(w)$ are accepted, then $\setminus(L(w))$ is disjoint from $\setminus(L)$.

Example

Assuming a DFA that reaches an accepting state on processing "1111001":

- $\setminus(L(1111001))$ is contained in $\setminus(L)$.

If it does not reach an accepting state:

- Further prefix testing is necessary to determine if $\setminus(L(1111001))$ is disjoint or neither.

Given the provided context, the exact transitions and states of the DFA are not detailed, but the approach remains consistent: simulate, analyze, and test prefixes to classify each $\setminus(L(w))$.