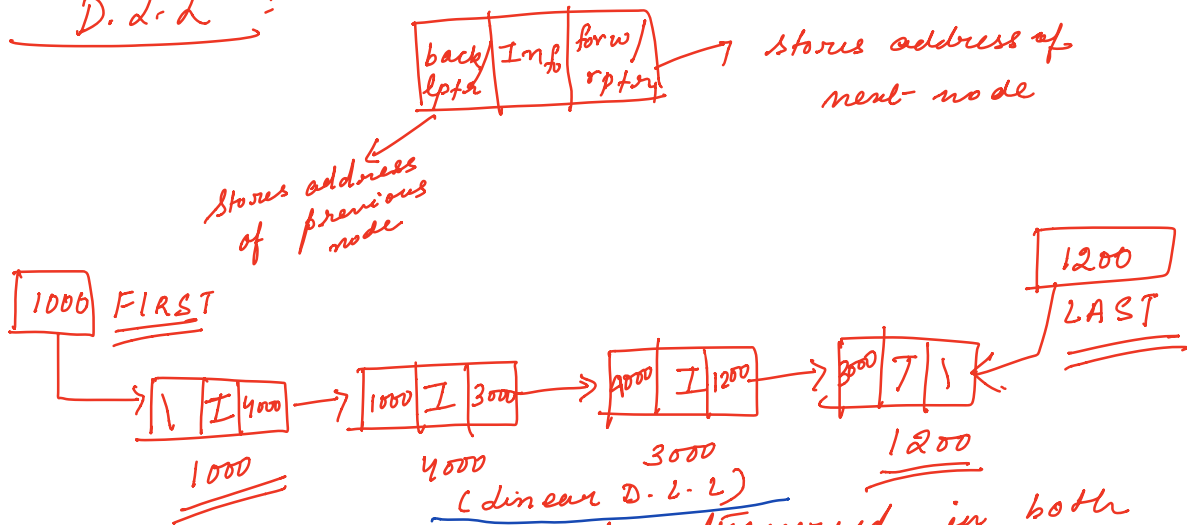


Dec-08/03/05/21

- Doubly d-list / 2-way list
- Applications of d-list
  - Poly. addition
  - Poly. multiplication
  - Linked dictionary
- Concept of garbage collection / compaction

D. d. d. :



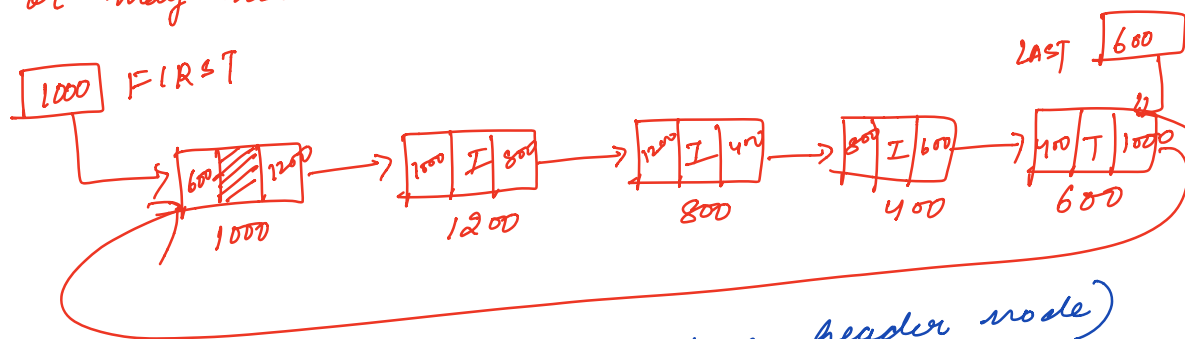
The 2-way list can be traversed in both directions → (i) forward direction (begin to end)  
(ii) backward direction (end to begin)

defn: A 2-way list is a linear collection of data elements called Nodes where each node  $N$  is divided into 3 parts:

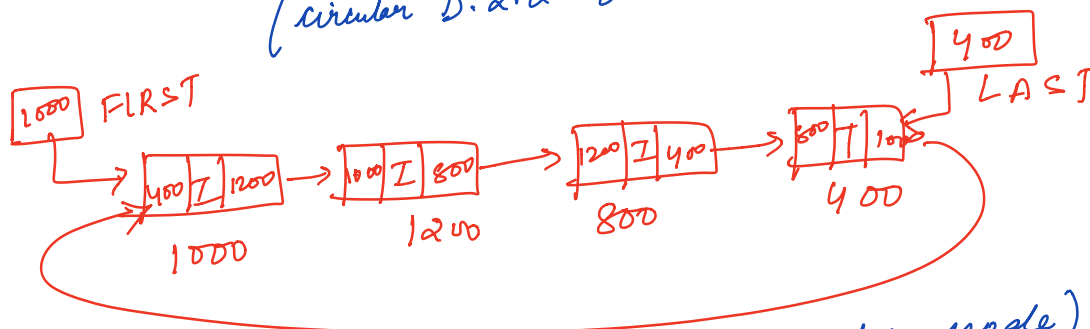
- An information field INFO contains user information.
- A pointer field FORW (rptr) which contains the address of next node in the d-list.
- A pointer field BACK (lptr) which contains the address of previous node in the d-list.

→ The list contains two list pointer variables FIRST and LAST which stores the address of first node and last node of D.L.L. resp.

\* D.L.L. may be either linear, circular may or may not contain a header node.



(circular D.L.L with a header node)



(circular D.L.L without a header node)

1) DL L- insert :

This algo. inserts a data item X to the left of a node whose address is M in a D.L.L. L & R represent the left most and right most node of D.L.L.

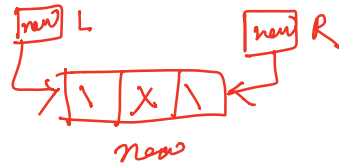
1) if avail = NULL, write "overflow" and return.

2) new = avail  
avail = link(avail)

3) set info(new) = X

4) if  $L = R = \text{NULL}$  then  
 $\text{LPTR}(\text{new}) = \text{RPTR}(\text{new}) = \text{NULL}$   
 set  $L = R = \text{new}$

// Insertion to an empty list



return

5) if  $M = L$  then

$\text{LPTR}(\text{new}) = \text{NULL}$  ✓

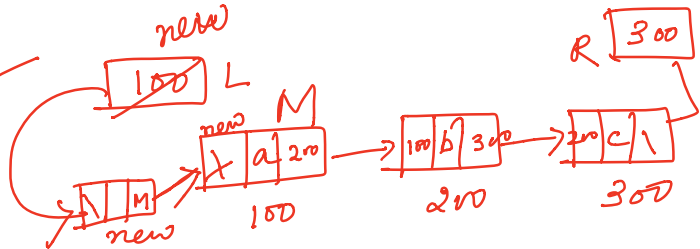
$\text{RPTR}(\text{new}) = M$

$\text{LPTR}(M) = \text{new}$

$L = \text{new}$

return

// Insert a node as left most node



6) //  $M$  is in middle. so insert new node left of  $M$  //

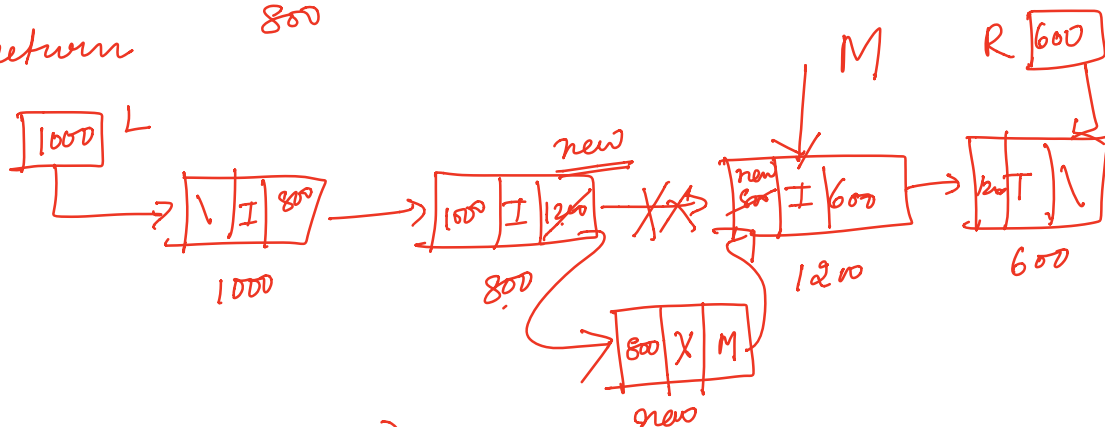
$\text{LPTR}(\text{new}) = \text{LPTR}(M)$  ✓

$\text{RPTR}(\text{new}) = M$

$\text{LPTR}(M) = \text{new}$

$\text{RPTR}(\text{LPTR}(\text{new})) = \text{new}$

return



2 D.L.L (delete):

This algo. deletes a node with an address doe from D.L.L with  $L$  and  $R$  as the pointer storing the address of left most and right most nodes resp.

Step 17 if  $L = R = \text{NULL}$ , write "underflow", return.

27 if  $L = R \neq \text{NULL}$  // single node before deletion  
then set  $L = R = \text{NULL}$

else if  $\text{LOC} = L$  // delete the left most node  
then  $L = \text{RPTR}(L)$   
 $\text{LPTR}(L) = \text{NULL}$

else if  $\text{LOC} = R$  // delete the right most node  
then  $R = \text{LPTR}(R)$   
 $\text{RPTR}(R) = \text{NULL}$

else // delete from middle

$\text{RPTR}(\text{LPTR}(\text{LOC})) = \text{RPTR}(\text{LOC})$   
 $\text{LPTR}(\text{RPTR}(\text{LOC})) = \text{LPTR}(\text{LOC})$

3. return.

