

# S O T N I R G

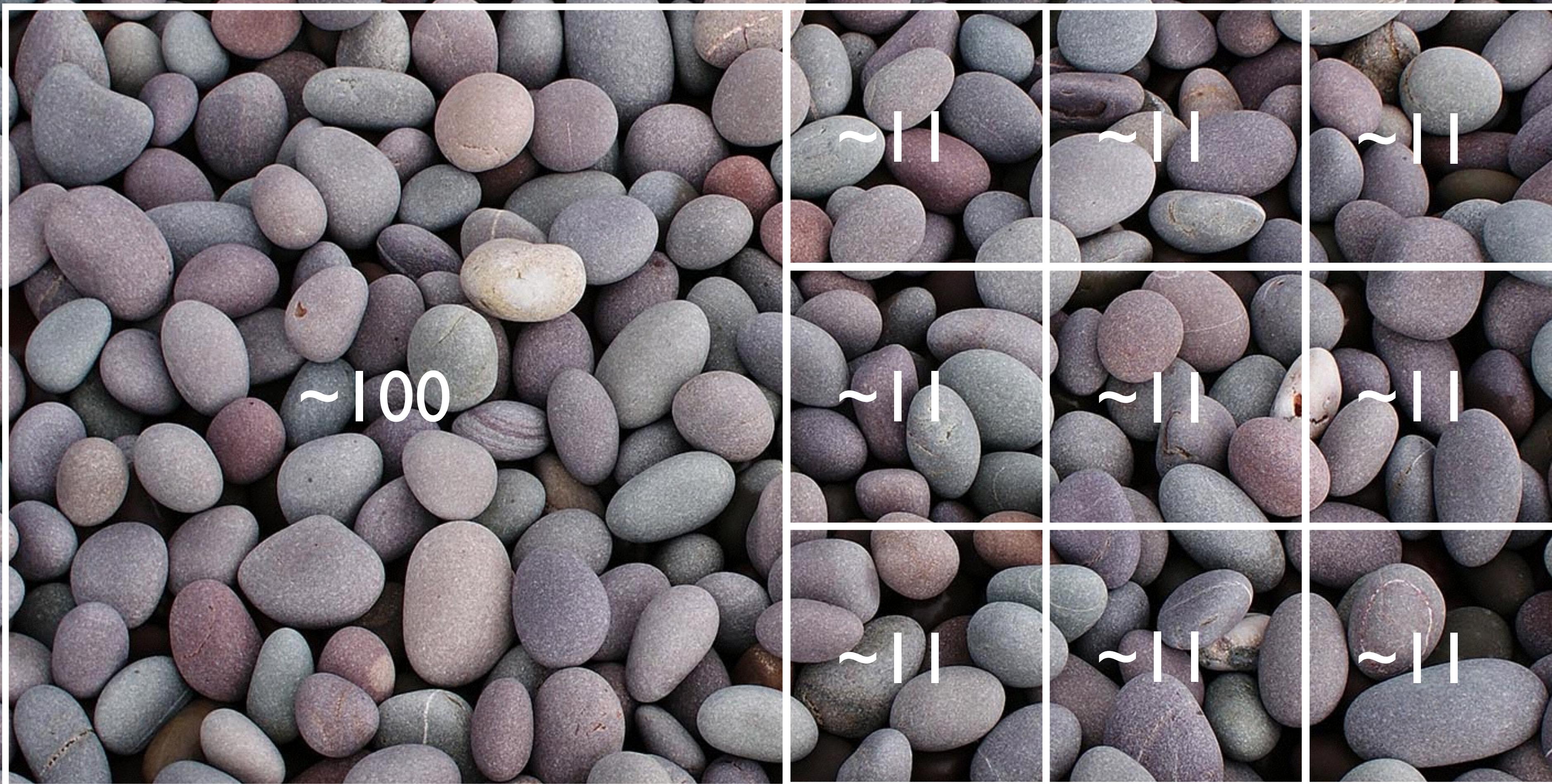


*Algorithms & Analysis*

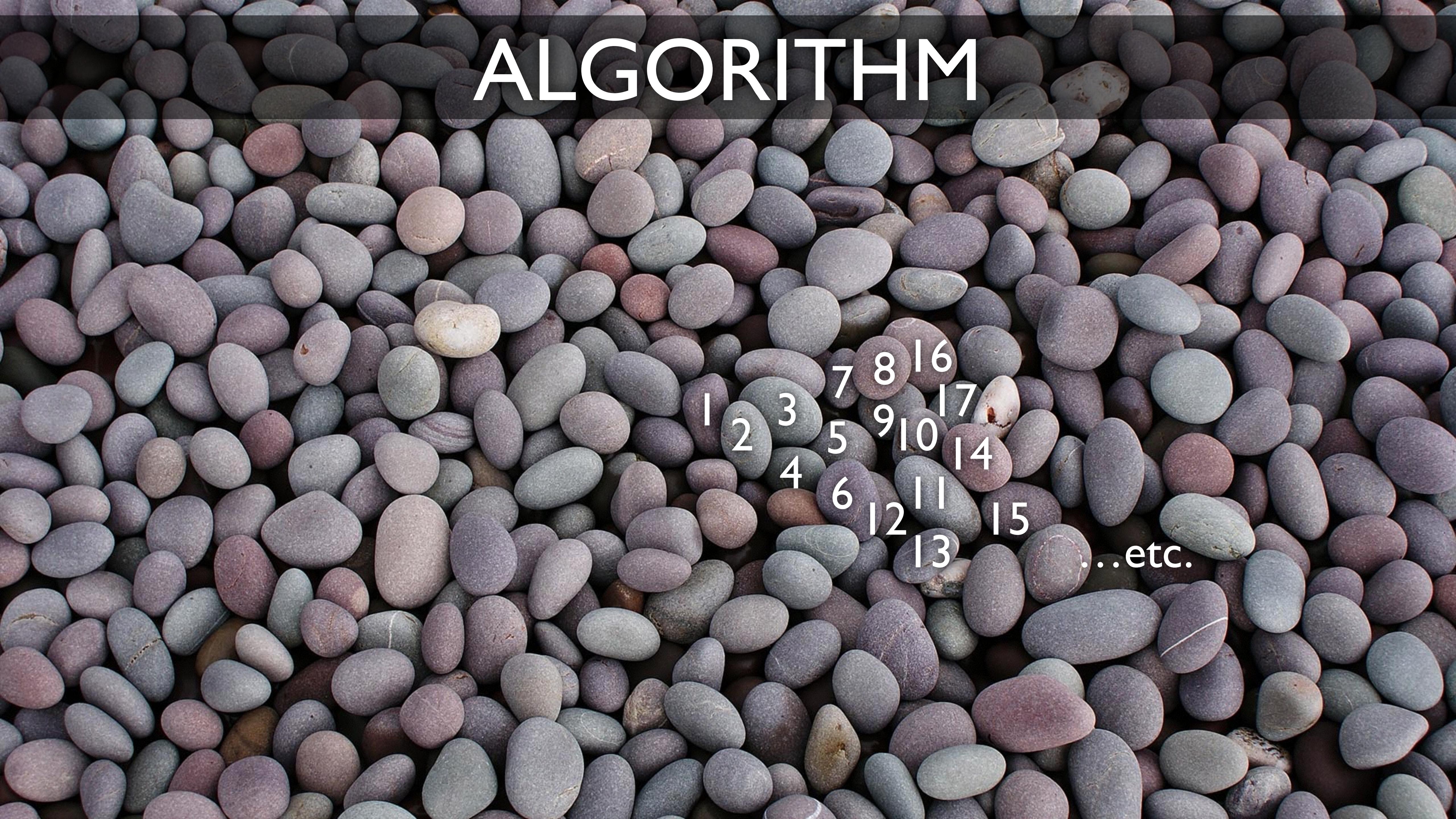
# But first: how many pebbles?



# HEURISTIC



# ALGORITHM

A close-up photograph of a large pile of dark, smooth, rounded stones or pebbles. The stones are various shades of grey, black, and brown, with some showing signs of wear and small white spots. They are densely packed, filling the entire frame.

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 ...etc.

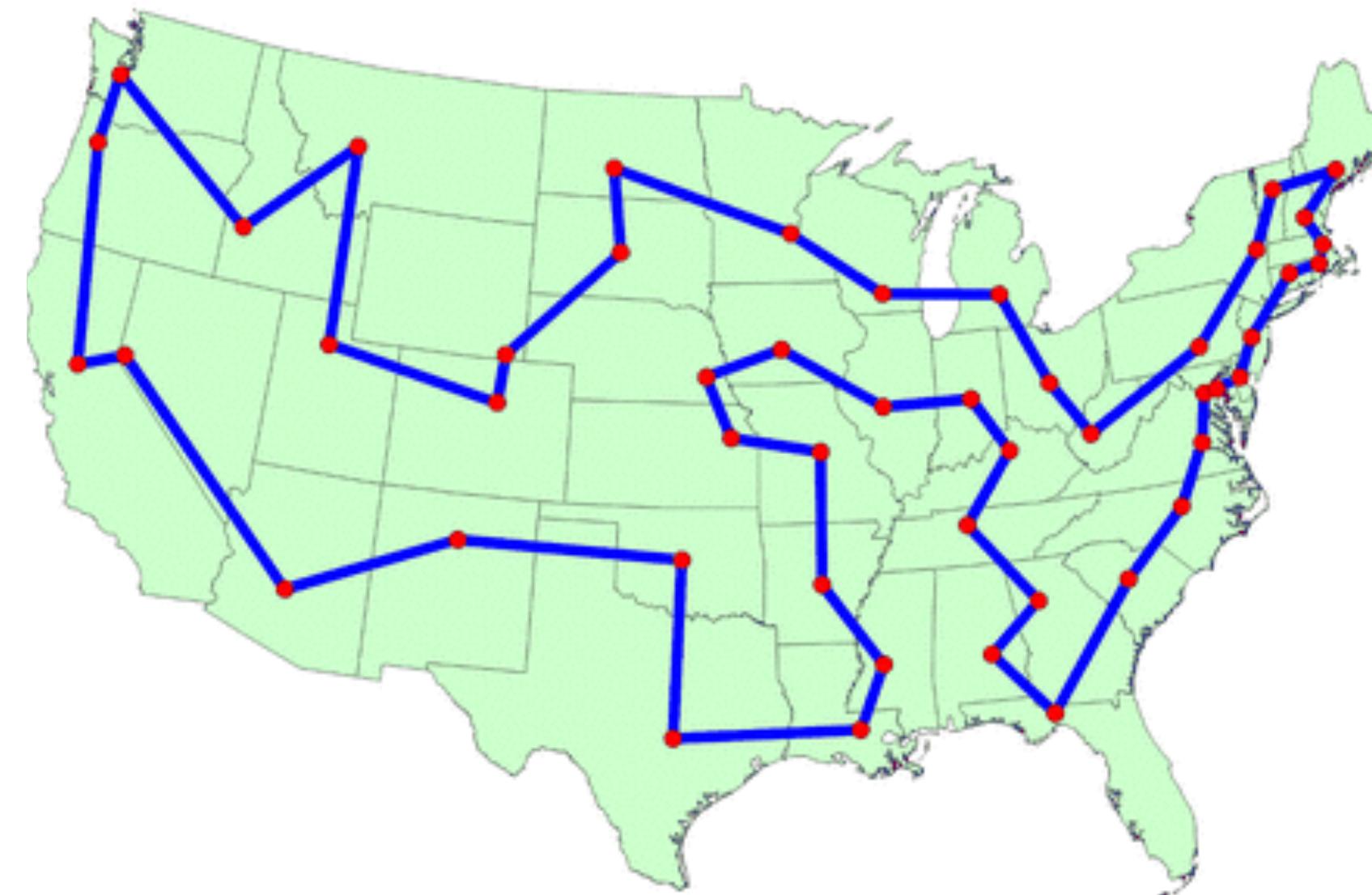
# Algorithms

- **Step-by-step instructions (deterministic)**
- **Complete** (gets you an answer)
- **Finite** (...given enough time)
- **Efficient** (doesn't waste time getting you the correct answer)
- **Correct** (the answer isn't just close, it is true)
- Downside: some problems are very **hard / slow**

*Often we loosely call functions algorithms, because much of the time a function is implementing an algorithm.*

# Heuristics

- Not necessarily correct (but gets you a "*good enough*" answer)
- Advantage: *fast* (often way faster than an algorithm)
- Famous example: the Traveling Salesman Problem



# How can we compare algorithms?

# THE *BIG*

A man in a dark tuxedo and bow tie looks up in awe at a massive, red and blue robot standing on a bridge. The robot has a large, segmented torso and arms, and a head with a prominent, bulbous nose. In the background, a city skyline is visible under a dark sky.

THE

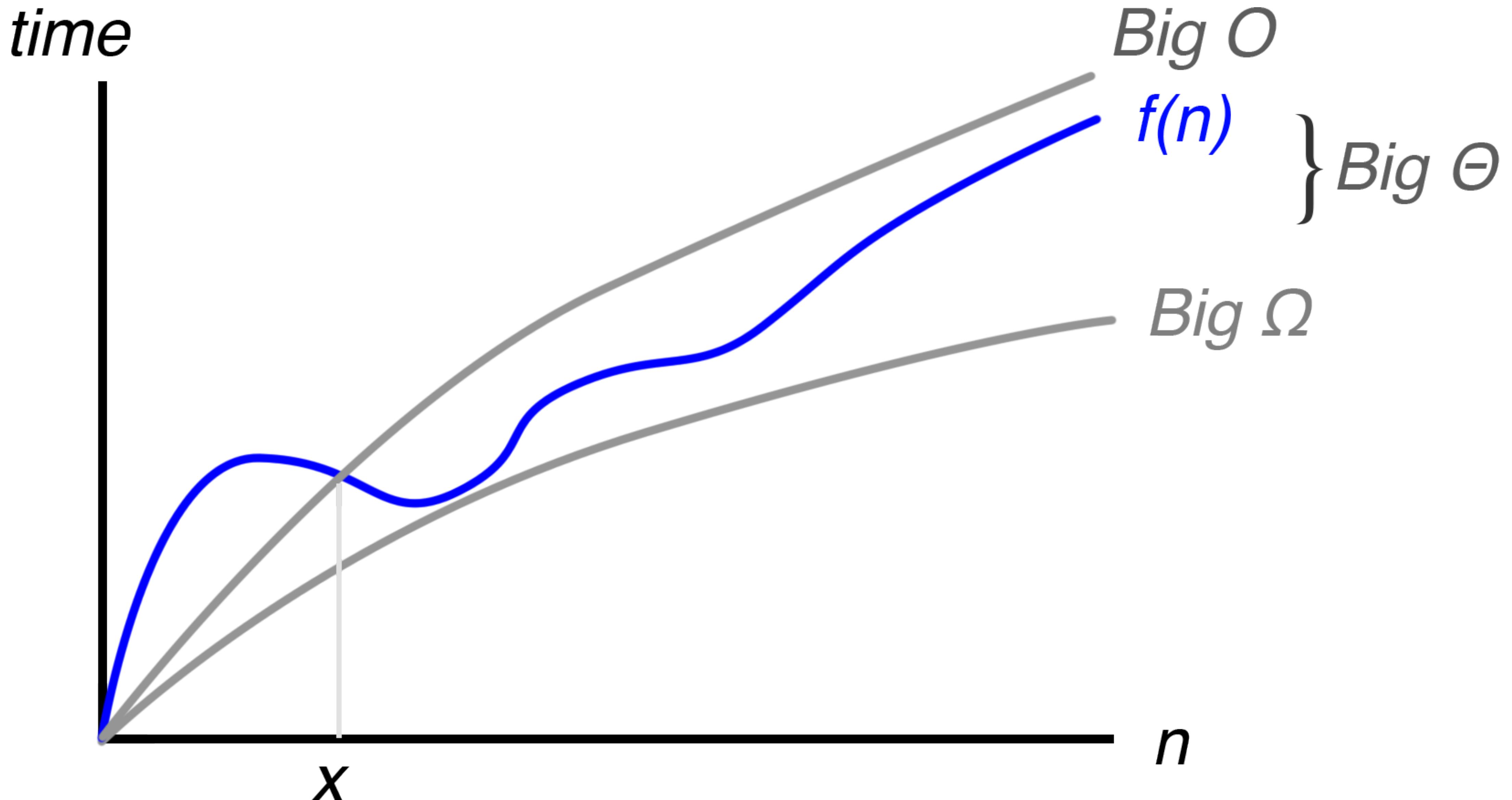
*BIG*

# Algorithm Analysis: Big O Notation

- A **comparative** way to classify different algorithms
- Based on **shape of growth curve** (**time vs input size(s)**)
- For **big enough** inputs
  - Might not be true when  $n$  is small, but who cares when  $n$  is small?
- Establishing an **upper bound** on the time
  - Not worse than this. Might be better, but it ain't worse!
- Including just the **highest order** term
  - In  $f(n) = n^3 + 5n + 3$ , only  $n^3$  matters as  $n$  gets large
- **Ignores constants** (mostly irrelevant;  $0.1 \cdot n^2$  will overtake  $10 \cdot n$ )



# What?



# Big O: comparative

- A very coarse, broad tool — big simplification
- Only useful when algorithms have *different* Big O notations
  - $O(n)$  will always beat  $O(n^2)$ , for *big* enough  $n$
- If two algorithms have the same Big O, we don't know much.  
One might actually be quite slower than the other.

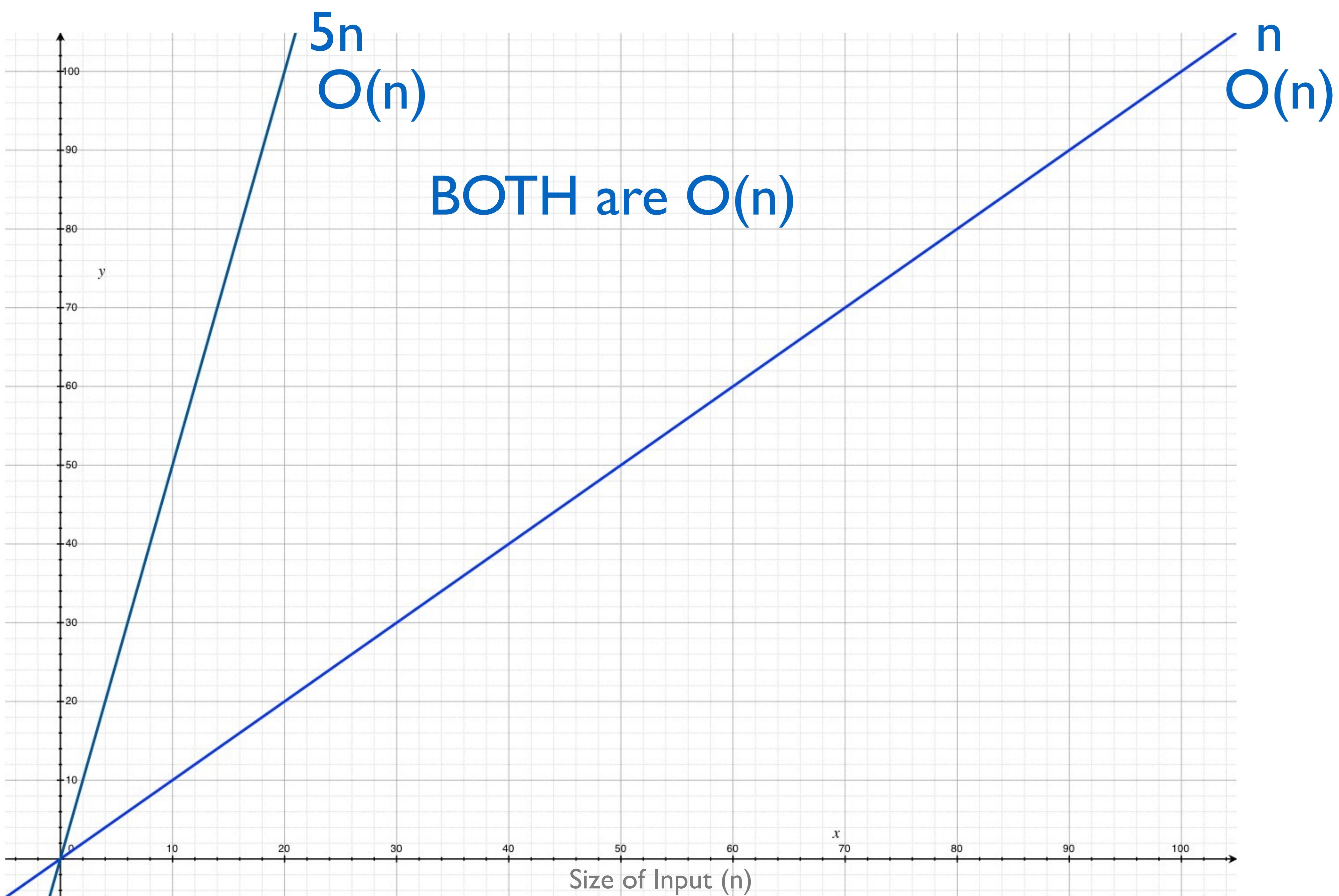


# Two Linear Functions

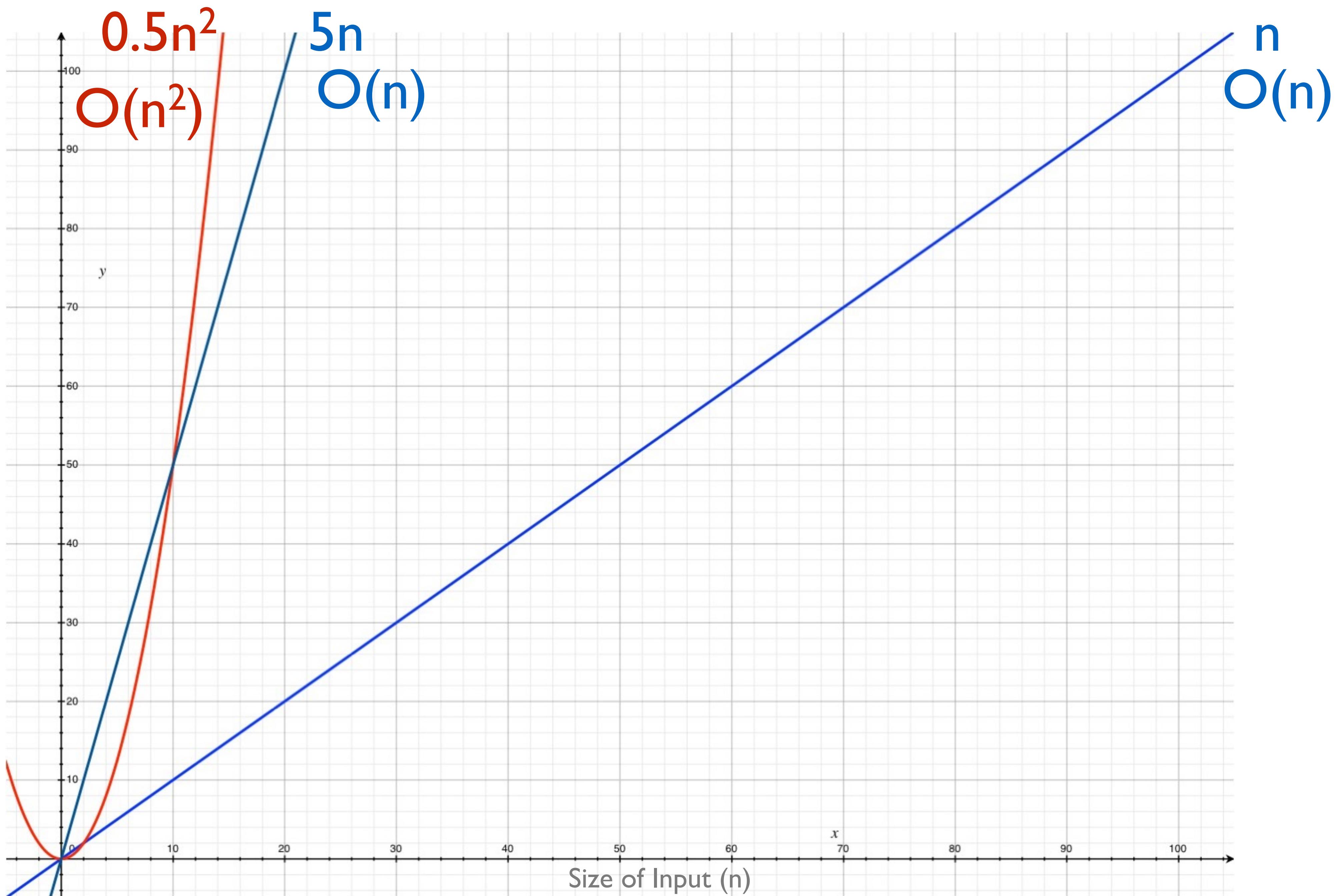
```
function findColors (arr) {  
  var colors = {  
    red: true,  
    orange: true,  
    yellow: true,  
    green: true,  
    blue: true  
  };  
  arr.forEach(function (val, i) {  
    if (colors[val]) console.log(i, val);  
  });  
}
```

```
function findColorsSlow (arr) {  
  arr.forEach(function (val, i) {  
    if (val === 'red') console.log(i, val);  
  });  
  arr.forEach(function (val, i) {  
    if (val === 'orange') console.log(i, val);  
  });  
  arr.forEach(function (val, i) {  
    if (val === 'yellow') console.log(i, val);  
  });  
  arr.forEach(function (val, i) {  
    if (val === 'green') console.log(i, val);  
  });  
  arr.forEach(function (val, i) {  
    if (val === 'blue') console.log(i, val);  
  });  
}
```

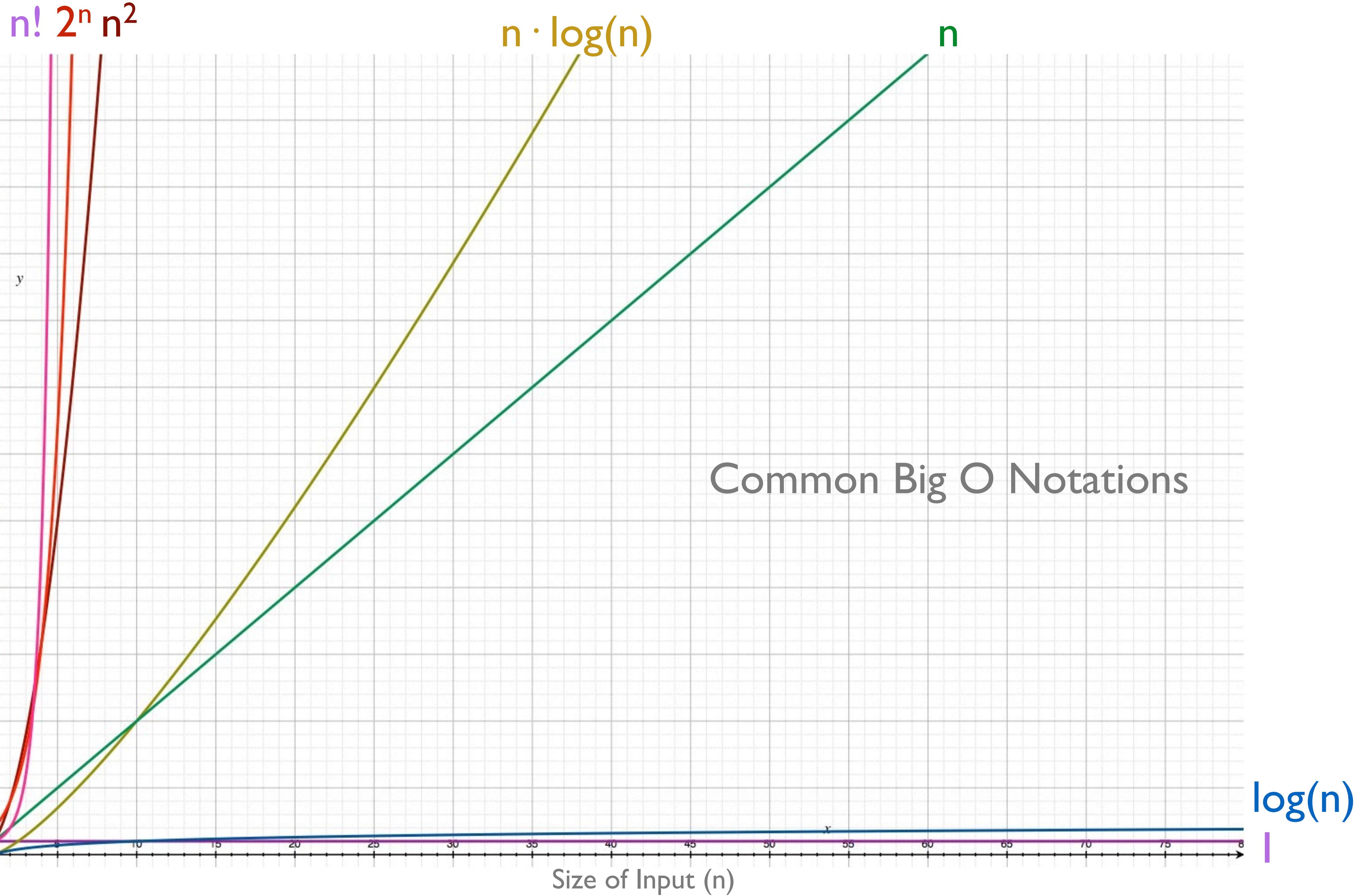
Time for  
Function  
to Complete



Time for  
Function  
to Complete



Time for  
Function  
to Complete





\*Source: Skiena, The Algorithm Design Manual

# Time Complexities (if 1 op = 1 ns)

input size n	$\log n$	n	$n \cdot \log n$	$n^2$	$2^n$	$n!$
10	0.003 μs	0.01 μs	0.03 μs	0.1 μs	1 μs	3.63 ms
20	0.004 μs	0.02 μs	0.09 μs	0.4 μs	1 ms	77.1 years
30	0.005 μs	0.03 μs	0.15 μs	0.9 μs	1 sec	8.4 × 10 <sup>15</sup> yrs
40	0.005 μs	0.04 μs	0.21 μs	1.6 μs	18.3 min	
50	0.006 μs	0.05 μs	0.28 μs	2.5 μs	13 days	
100	0.007 μs	0.10 μs	0.64 μs	10.0 μs	4 × 10 <sup>13</sup> yrs	
1 000	0.010 μs	1.00 μs	9.97 μs	1 ms		
10 000	0.013 μs	10.00 μs	~130.00 μs	100 ms		
100 000	0.017 μs	100.00 μs	1.7 ms	10 sec		
1 000 000	0.020 μs	1 ms	19.9 ms	16.7 min		
10 000 000	0.023 μs	10 ms	230.0 ms	1.16 days		
100 000 000	0.027 μs	100 ms	2.66 sec	115.7 days		
1 000 000 000	0.030 μs	1 sec	29.90 sec	31.7 years		



# Time Complexities

Big O	Name	Think	Example
$O(1)$	<i>Constant</i>	Doesn't depend on input	get array value by index
$O(\log n)$	<i>Logarithmic</i>	Using a tree	find min element of BST
$O(n)$	<i>Linear</i>	Checking (up to) all elements	search through linked list
$O(n \cdot \log n)$	<i>Loglinear</i>	A tree for each element	merge sort average & worst case
$O(n^2)$	<i>Quadratic</i>	Checking pairs of elements	bubble sort average & worst case
$O(2^n)$	<i>Exponential</i>	Generating all subsets	brute-force n-long binary number
$O(n!)$	<i>Factorial</i>	Generating all permutations	the Traveling Salesman



Data Structure	Time Complexity								
	Average				Worst				
	Access	Search	Insertion	Deletion		Access	Search	Insertion	Deletion
Array	0(1)	0(n)	0(n)	0(n)	0(1)	0(n)	0(n)	0(n)	0(n)
Stack	0(n)	0(n)	0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(1)
Singly-Linked List	0(n)	0(n)	0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(1)
Doubly-Linked List	0(n)	0(n)	0(1)	0(1)	0(n)	0(n)	0(1)	0(1)	0(1)
Skip List	0(log(n))	0(log(n))	0(log(n))	0(log(n))	0(n)	0(n)	0(n)	0(n)	0(n)
Hash Table	-	0(1)	0(1)	0(1)	-	0(n)	0(n)	0(n)	0(n)
Binary Search Tree	0(log(n))	0(log(n))	0(log(n))	0(log(n))	0(n)	0(n)	0(n)	0(n)	0(n)

*“By understanding sorting, we obtain an amazing amount of power to solve other problems.”*

– STEVEN SKIENA, THE ALGORITHM DESIGN MANUAL

# (Some) Classic Sorting Algorithms

- **Bubble**
- **Selection**
- **Insertion**
- **Merge: 1945 Jon von Neumann**
- **Quick: 1959 Tony Hoare**
- **Heap: 1964 J. W. J. Williams**
- **Radix: 1887 Hermann Hollerith, for his Tabulating Machine**
- **Bogo?**

# Bubble Sort

6 5 3 1 8 7 2 4

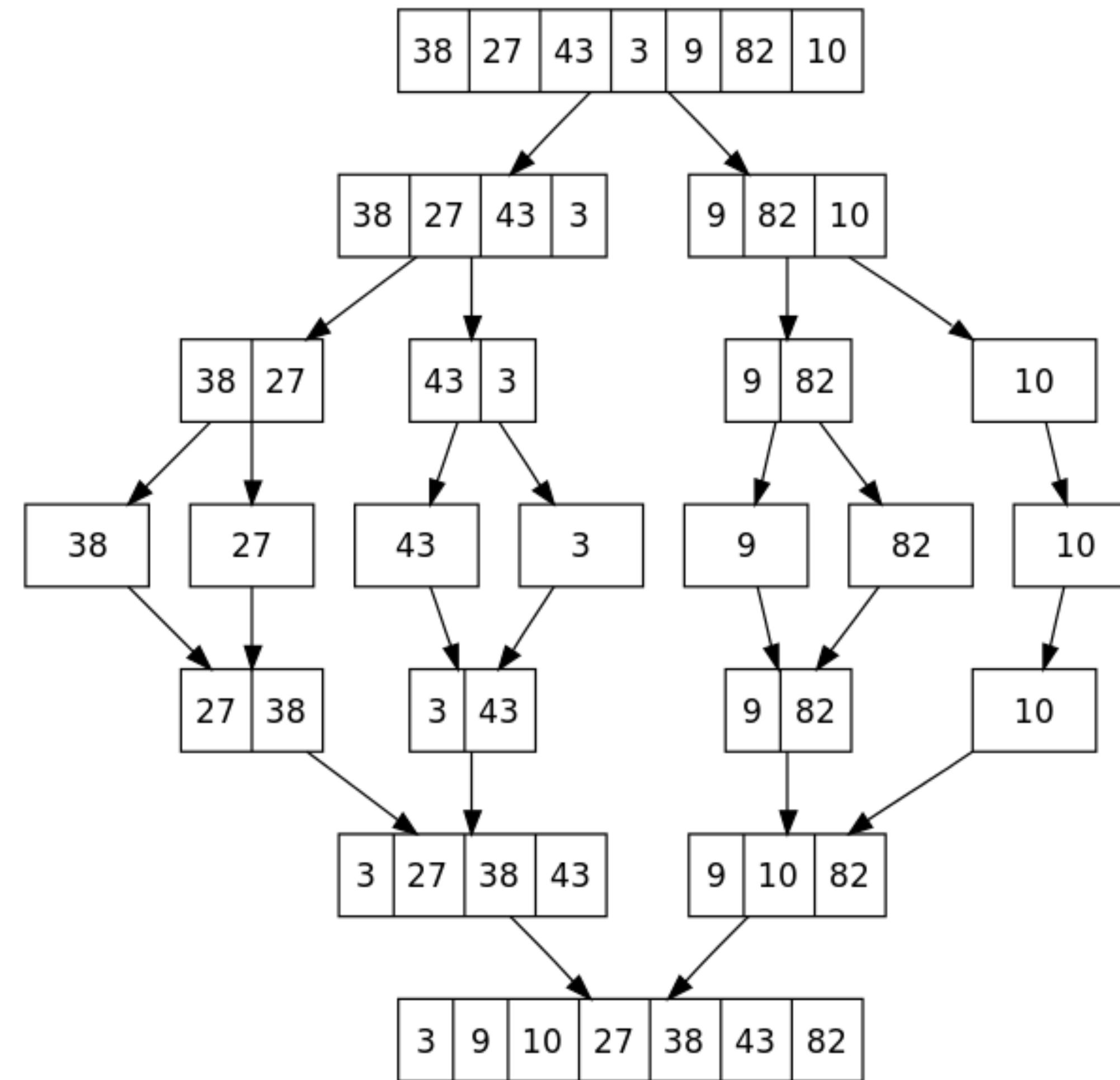
# Bubble Sort

1. Loop over elements
2. Swap anything that's out of order
3. Repeat 1-2 until there are no swaps

# Merge Sort

6 5 3 1 8 7 2 4

# Merge Sort



# Merge Sort (iterative)

1. Divide array of  $n$  elements into  $n$  arrays of 1 element
2. Merge neighboring arrays in sorted order
3. Repeat 2 until there's only one array

# Merge Sort (recursive)

1. If array is one element, good job it's sorted!
2. Otherwise, split the array and merge sort each half
3. Merge combined halves into sorted whole

# Big O

	Bubble Sort	Merge Sort
Time	$O(n^2)$	$O(n \cdot \log n)$
Space	$O(1)$	$O(n)$

# Why is merge sort faster?



# Merge Sort Speedup

- Combining two lists that are each already sorted into one list that is sorted is a linear time operation
- There are  $\log_2(n)$  steps needed to go from  $n$  lists of one item each to one list of  $n$  items

# Special Note

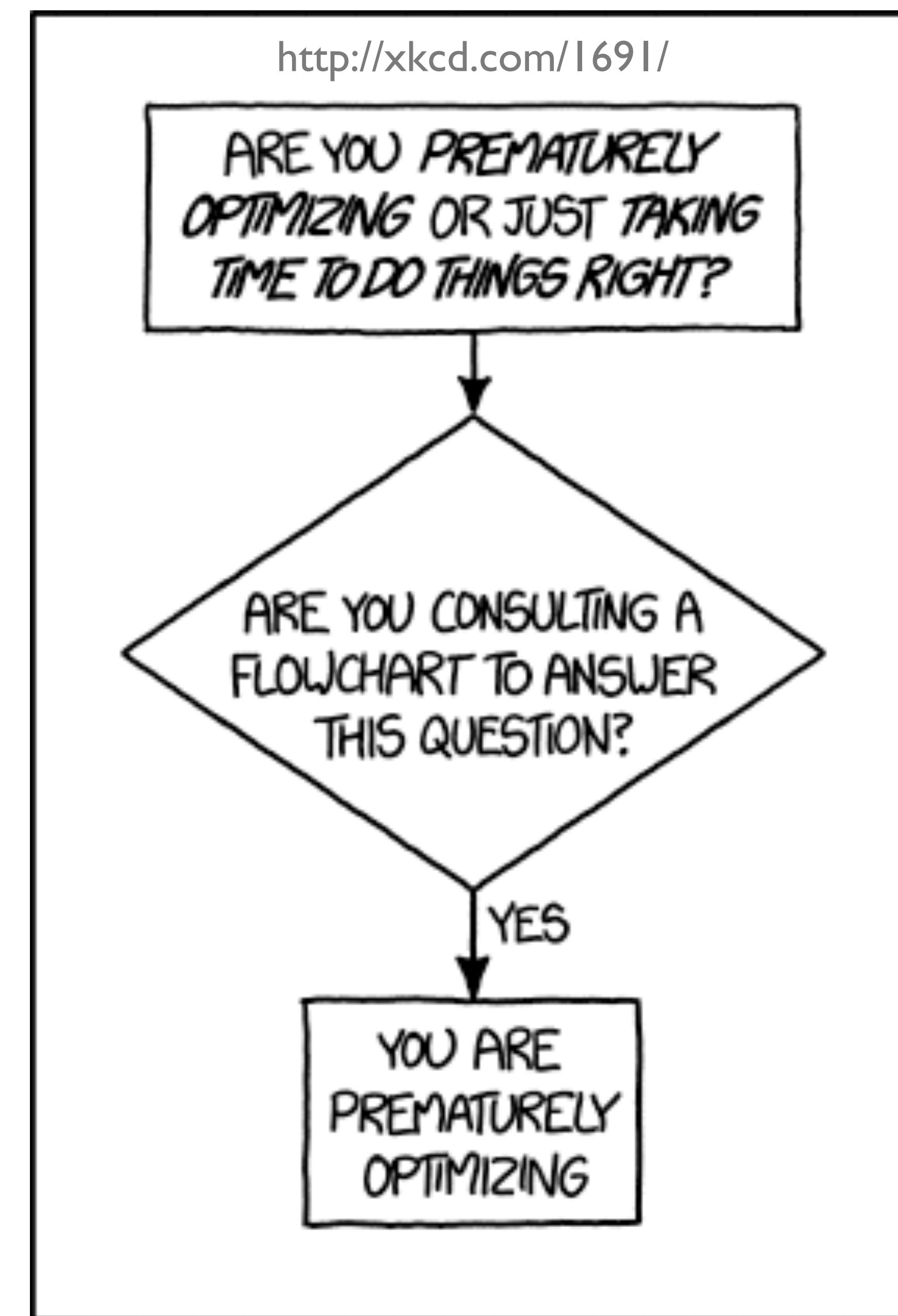


# Rob Pike's 5 Rules of Programming

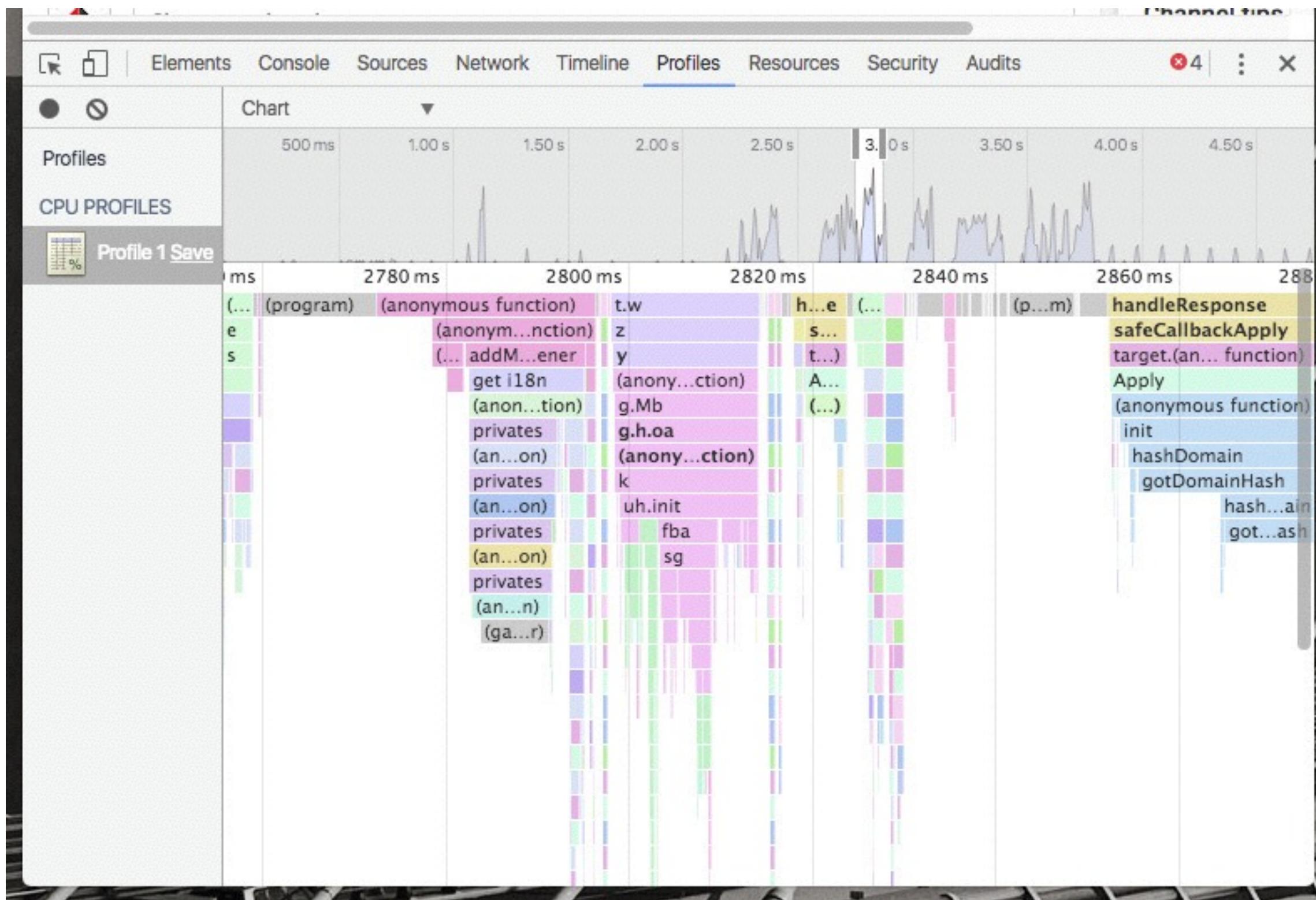
Bell Labs  
Unix Team  
UTF-8  
Go Language  
...and a lot more

# 1

- You can't tell where a program is going to spend its time. Bottlenecks occur in surprising places, so don't try to second guess and put in a speed hack until you've proven that's where the bottleneck is.



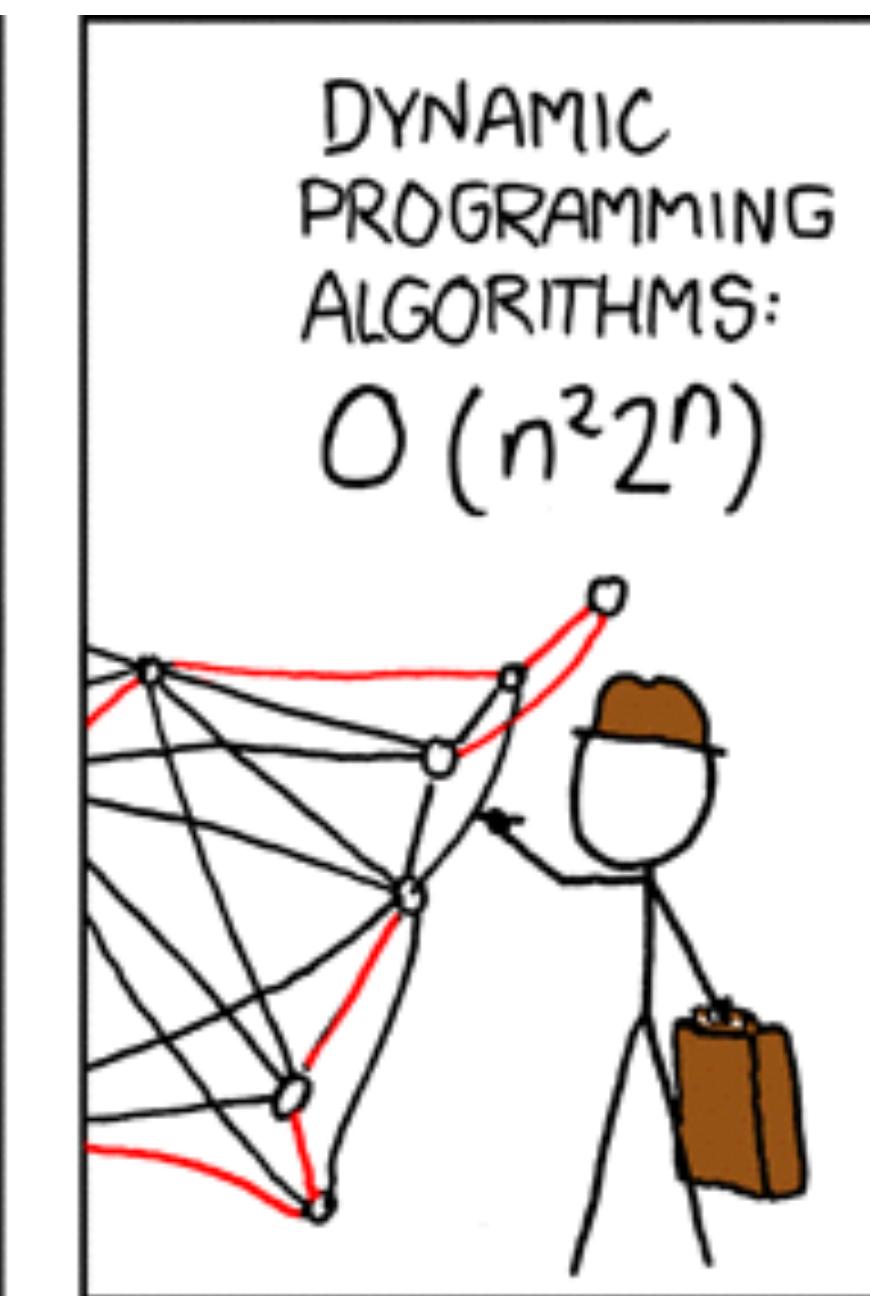
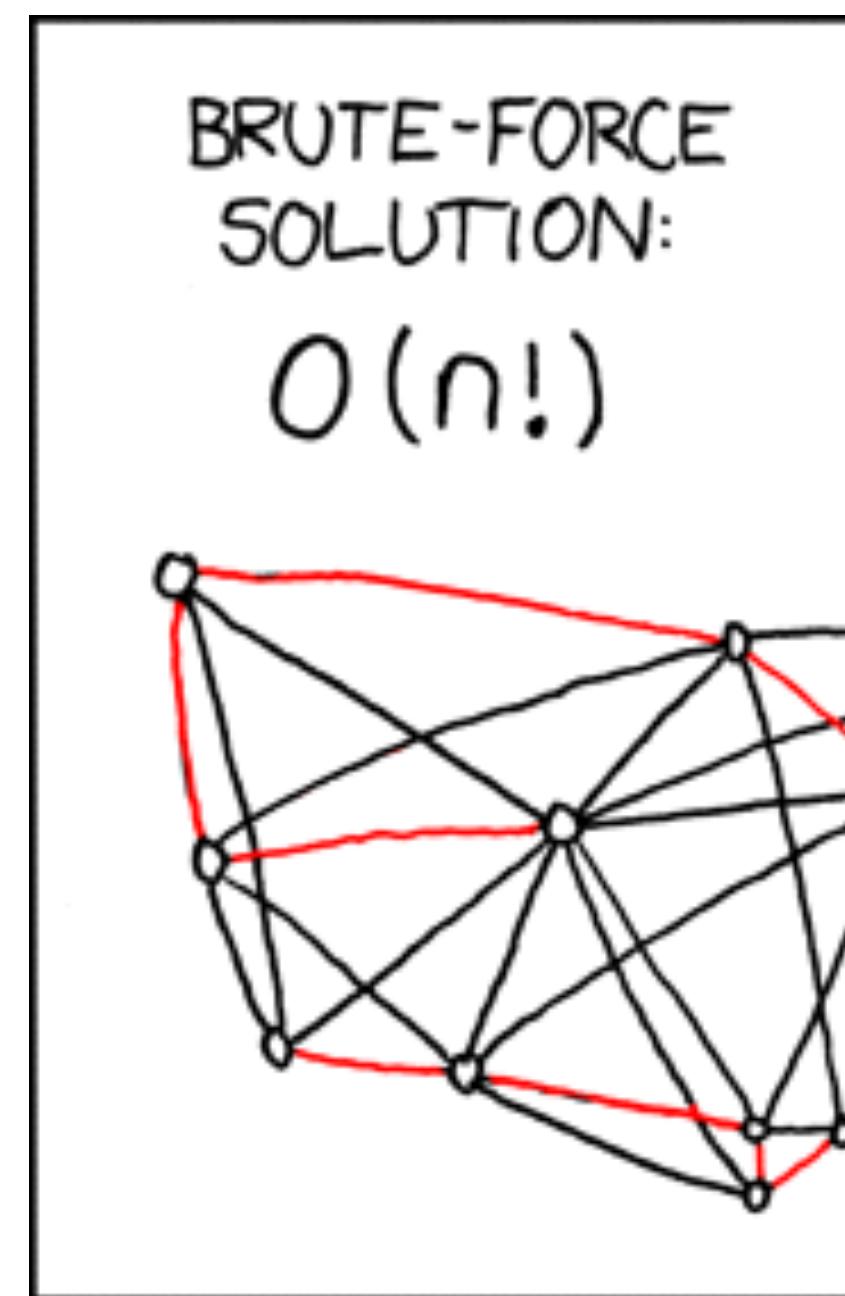
# 2



- Measure. Don't tune for speed until you've measured, and even then don't unless one part of the code overwhelms the rest.

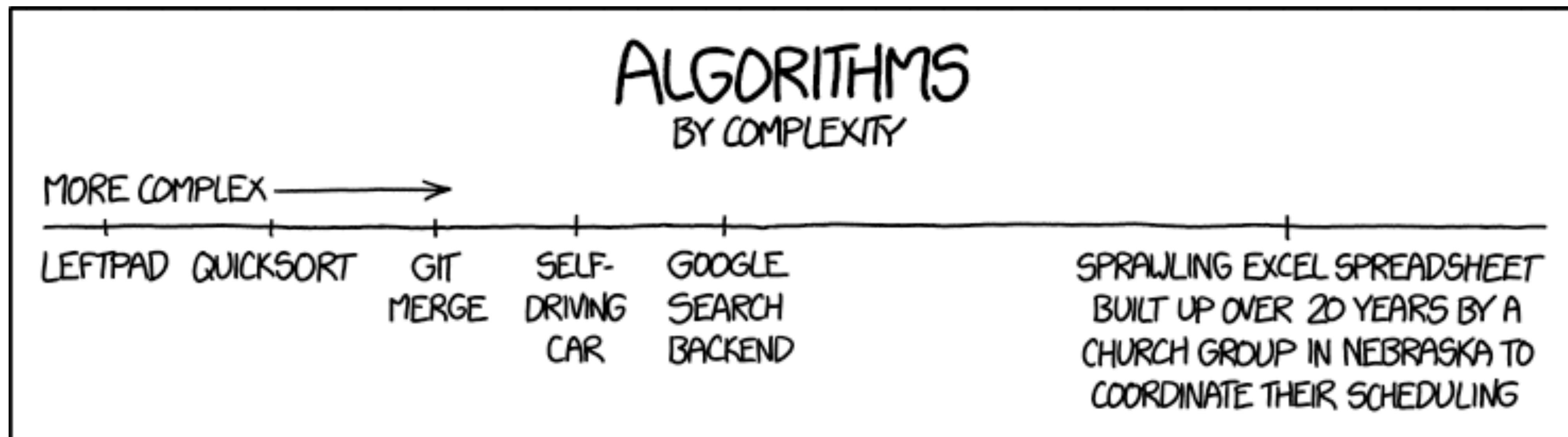
# 3

- Fancy algorithms are slow when  $n$  is small, and  $n$  is usually small. Fancy algorithms have big constants. Until you know that  $n$  is frequently going to be big, don't get fancy.



# 4

- Fancy algorithms are buggier than simple ones, and they're much harder to implement. Use simple algorithms as well as simple data structures.



# 5

- ◎ Data dominates. If you've chosen the right data structures and organized things well, the algorithms will almost always be self-evident. Data structures, not algorithms, are central to programming.

# WORKSHOP