

Algorithms and their complexity

Lecture 9a

Waterford Institute of Technology

March 9, 2015

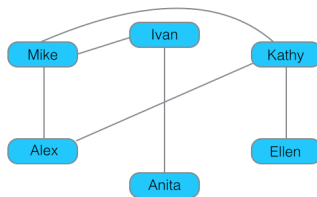
John Fitzgerald

Algorithm

Description

What is a software algorithm?

- Program that solves problem
- Algorithm examples:
 - Sorting data
 - Searching data
 - Graphs (social networks)
 - Cryptography
 - Image processing
 - From labs:
 - Generate PIN
 - Validate PIN



Modelling social network

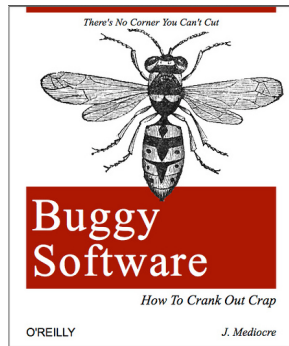
Algorithm output

Expect correct result for all valid inputs

Easier said than done

Examples of incorrect output:

- Timsort
 - Extensively used on Android
 - Bug recently discovered
- Binary search
 - Small & conceptually simple
 - Yet took years to eliminate all bugs for all inputs



Algorithm performance

How long to run?

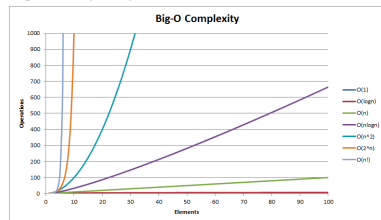
Goal not measurement run time
because of variations in

- Computer speeds
- Language implementation
- Inputs

Instead, determine inherent
complexity

- Classes of complexity exist
- Provide comparison of
growth of number
computing operations

Big-O Complexity Chart



Linear search algorithm

How long to run?

Consider the linear search algorithm below

Running times could vary significantly:

- If *String search* near start of *String[] target*
 - method returns immediately
- But if towards end of list
 - return might be much later
 - providing very different run times for same algorithm

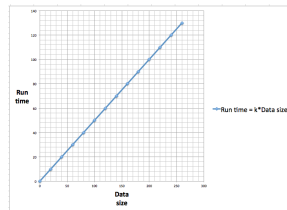
```
boolean linearSearch(String search, String[] target)
{
    for (int i = 0; i < target.length; i += 1)
        if (target[i].contains(search))
            return true;
    return false;
}
```

Algorithm efficiency

Choices available

We have a choice to make:

- Best case
 - Considering all possible inputs
 - Linear search runs in constant time
 - Returns immediately
- Worst case
 - Search time proportional input size
 - Linear search is *linear* in size list
- Average



We will choose worst case.

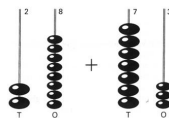
- We will seek an upper bound on number operations in algorithm.

Counting operations

Agree some rules

Basic assumptions

- Concept of random access machine (RAM)
- An abstract computer
- Measure size of input
- Sequential execution operations
- Operations take constant time
 - Assignment ($a = 10$)
 - Comparison ($a == 10$)
 - Arithmetic ($a/10$)
 - Memory access ($b[i]$)



Counting operations

Simple example

Approximate number operations or steps in *countOperations*:

- $2001 + 2n + 2n^2$

```
static int countOperations(int n)
{
    int ans = 0;
    for (int i = 0; i < 1000; i += 1)
        ans += 1;
    for (int i = 0; i < n; i += 1)
        ans += 1;
    for (int i = 0; i < n; i += 1)
        for (int j = 0; j < n; j += 1)
            ans += 1;
    return ans;
}
```

Number
operations

1

2000

$2n$

$2n^2$

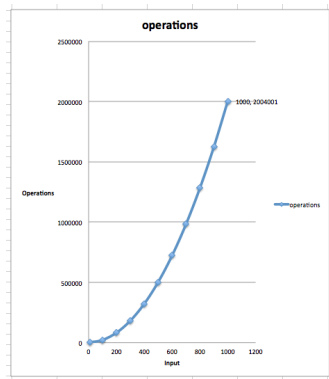
Count operations

Number steps or operations

$2001 + 2n + 2n^2$ (ignoring *for* overhead)

- If n small output dominated by 2000 term
- For large n 2000 term insignificant
 - If n is 10 output is 2221
 - If n is 10,000 output exceeds 200 million ($2.0002E+08$)

```
static int countOperations(int n)
{
    int ans = 0;
    for (int i = 0; i < 1000; i += 1)
        ans += 1;
    for (int i = 0; i < n; i += 1)
        ans += 1;
    for (int i = 0; i < n; i += 1)
        for (int j = 0; j < n; j += 1)
            ans += 1;
    return ans;
}
```

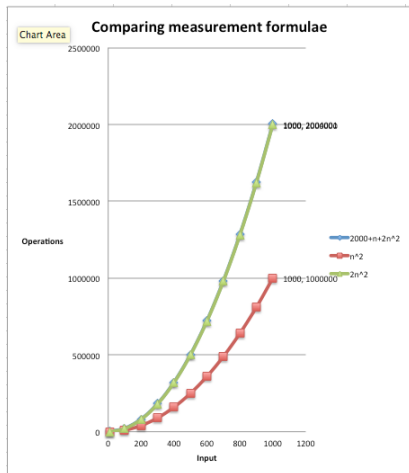


Algorithm complexity

Comparing measurement formulae

In the case of $2001 + 2n + 2n^2$
we use n^2

- We use the highest order variable
- And disregard any coefficient
- Provides approximate upper bound measurement
- Importantly, it provides representation of growth of complexity as input becomes very large.



Algorithm complexity

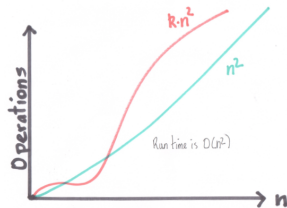
Use of asymptotic notation

Last example: Some constant times n^2 provides upper bound on number operations

Asymptotic notation used to describe growth of algorithm operations as input size approaches infinity.

We now introduce **Big O** notation

- Running time of $2001 + 2n + 2n^2$
 - $O(n^2)$ or
 - Big- $O(n^2)$**

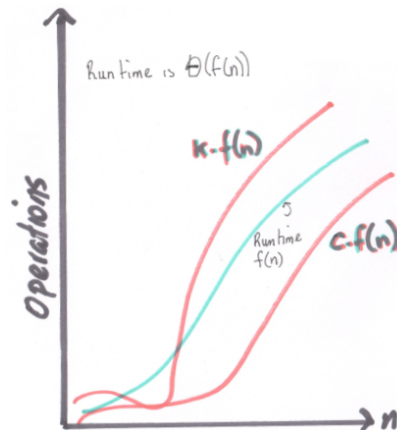


Algorithm complexity

Further asymptotic notation

Big Θ (Big Theta)

- For large n :
 - $k \cdot f(n)$ provides upper bound
 - $c \cdot f(n)$ provide lower bound
- k and c are some constants.

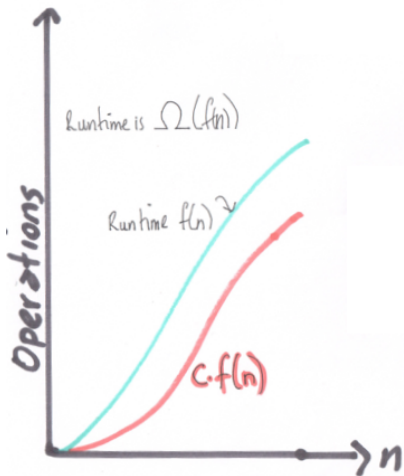


Algorithm complexity

Further asymptotic notation

Big Ω (Big Omega)

- For large n , $c \cdot f(n)$ provides lower bound run time
- c is some constant.

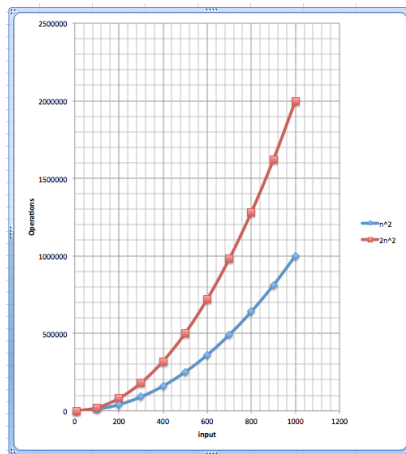


Algorithmic complexity

Alternative measurement approach

Tilde notation

- Proposed by Robert Sedegwick
- Denoted by: $\sim f(n)$
- Differs from Big-O
 - Coefficient of term with highest exponent retained
- Consider: $f(x) = 2000 + 2n + 2n^2$
 - $\sim 2n^2$ is considered the complexity of this algorithm



Big-O complexity

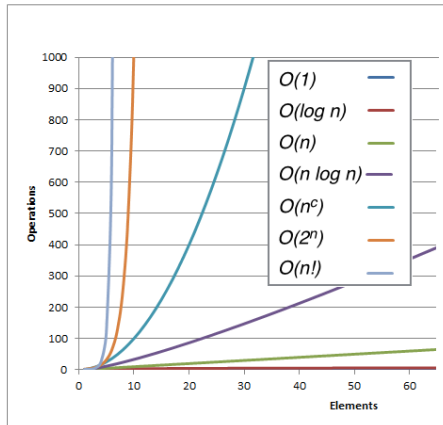
Categories

- Constant $O(1)$: independent of input
- Linear $O(n)$: find smallest item unsorted array
- Logarithmic $O(\log n)$: binary search
- Linearithmic $O(n \log n)$: mergesort
- Polynomial $O(n^c)$, $c > 0$: selection sort quadratic ($c == 2$)
- Exponential $O(2^n)$: Binary exponential backoff (networking)
- Factorial $O(n!)$: Brute force search travelling salesman problem

Categories

- Constant $O(1)$
- Logarithmic $O(\log n)$
- Linear $O(n)$
- Linearithmic $O(n \log n)$
- Polynomial $O(n^c)$
- Exponential $O(2^n)$
- Factorial $O(n!)$

Big-O Complexity



Algorithm complexity

Acceptable categories

Very large performance differences between categories

Does this really matter?

- For small problems, not really
- For very large problems, definitely yes.

Size array	Merge sort	Selection sort
20,000	2ms	868ms
100,000	10ms	13s 230ms
1,000,000	117ms	21m 55s 254ms
500,000,000	2m 2s 652ms	10.25 years

Search algorithms

Binary search

Linear search complexity is $O(n)$.

A more efficient *Binary search* has complexity $O(\log n)$

- Works with a sorted array
- Array considered as upper and lower half
- Check carried out: is item in upper or lower?
- Check repeated in the half array containing item.
- This search method repeated until item found if it exists.
- Each iteration reduces the search space in two.
- This explains the $O(\log n)$ order of complexity.

Search algorithms

Binary search

index ->	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]
iteration	{11,	12,	14,	17,	19,	20,	22,	24}
1	lo			mid				hi
2					lo	mid		hi
3						mid lo		hi

Binary search (22)

Search algorithms

Binary search

index ->	[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]
iteration	{11, 12, 14, 17, 19, 20, 22, 24}							
1	lo			mid				hi
2	lo	mid	hi					
3	hi mid lo							

Binary search (0)

Summary

- Methods to measure complexity
 - Elapsed time
 - Big O
 - Big Θ
 - Big Ω
 - Tilde notation
- Example algorithms
 - Selection sort
 - Merge sort
 - Linear search
 - Binary search
- Relative performance of algorithms

Referenced Material

1. Algorithms (Khan Academy)

www.khanacademy.org/computing/computer-science/algorithms/asymptotic-notation/a/big-o-notation

[Accessed 2015-03-03]

2. TimSort bug fixed with formal methods <http://www.cwi.nl/news/2015/java-bug-fixed-formal-methods-cwi>

[Accessed 2015-03-03]

3. edX | MITx: 6.00.1x Introduction to Computer Science and Programming

<https://www.edx.org> [Accessed 2015-03-05]

Referenced Material

4. Robert Sedgewick : Algorithms for the masses

<http://osric.com/chris/accidental-developer/2012/04/robert-sedgewick-algorithms-for-the-masses/> [Accessed 2015-03-05]