

SPIR-V Specification

John Kessenich, Google and Boaz Ouriel, Intel

Version 1.00, Revision 4

March 11, 2016



Copyright © 2014-2016 The Khronos Group Inc. All Rights Reserved.

This specification is protected by copyright laws and contains material proprietary to the Khronos Group, Inc. It or any components may not be reproduced, republished, distributed, transmitted, displayed, broadcast, or otherwise exploited in any manner without the express prior written permission of Khronos Group. You may use this specification for implementing the functionality therein, without altering or removing any trademark, copyright or other notice from the specification, but the receipt or possession of this specification does not convey any rights to reproduce, disclose, or distribute its contents, or to manufacture, use, or sell anything that it may describe, in whole or in part.

Khronos Group grants express permission to any current Promoter, Contributor or Adopter member of Khronos to copy and redistribute UNMODIFIED versions of this specification in any fashion, provided that NO CHARGE is made for the specification and the latest available update of the specification for any version of the API is used whenever possible. Such distributed specification may be reformatted AS LONG AS the contents of the specification are not changed in any way. The specification may be incorporated into a product that is sold as long as such product includes significant independent work developed by the seller. A link to the current version of this specification on the Khronos Group website should be included whenever possible with specification distributions.

Khronos Group makes no, and expressly disclaims any, representations or warranties, express or implied, regarding this specification, including, without limitation, any implied warranties of merchantability or fitness for a particular purpose or non-infringement of any intellectual property. Khronos Group makes no, and expressly disclaims any, warranties, express or implied, regarding the correctness, accuracy, completeness, timeliness, and reliability of the specification. Under no circumstances will the Khronos Group, or any of its Promoters, Contributors or Members or their respective partners, officers, directors, employees, agents, or representatives be liable for any damages, whether direct, indirect, special or consequential damages for lost revenues, lost profits, or otherwise, arising from or in connection with these materials.

Khronos, SYCL, SPIR, WebGL, EGL, COLLADA, StreamInput, OpenVX, OpenKCam, glTF, OpenKODE, OpenVG, OpenWF, OpenSL ES, OpenMAX, OpenMAX AL, OpenMAX IL and OpenMAX DL are trademarks and WebCL is a certification mark of the Khronos Group Inc. OpenCL is a trademark of Apple Inc. and OpenGL and OpenML are registered trademarks and the OpenGL ES and OpenGL SC logos are trademarks of Silicon Graphics International used under license by Khronos. All other product names, trademarks, and/or company names are used solely for identification and belong to their respective owners.

Contents

1	Intro	oduction	8
	1.1	Goals	8
	1.2	About this document	8
	1.3	Extendability	9
	1.4	Debuggability	9
	1.5	Design Principles	9
	1.6	Static Single Assignment (SSA)	10
	1.7	Built-In Variables	10
	1.8	Specialization	10
	1.9	Example	11
2	Spec	ification	14
	2.1	Language Capabilities	14
	2.2	Terms	14
		2.2.1 Instructions	14
		2.2.2 Types	15
		2.2.3 Module	16
		2.2.4 Control Flow	16
	2.3	Physical Layout of a SPIR-V Module and Instruction	18
	2.4	Logical Layout of a Module	19
	2.5	Instructions	20
		2.5.1 SSA Form	20
	2.6	Entry Point and Execution Model	21
	2.7	Execution Modes	21
	2.8	Types and Variables	21
	2.9	Function Calling	22
	2.10	Extended Instruction Sets	22
	2.11	Structured Control Flow	23
	2.12	Specialization	24
	2.13	Linkage	25
	2.14	Relaxed Precision	25
	2.15	Debug Information	26
		2.15.1 Function-Name Mangling	26
	2.16	Validation Rules	27
		2.16.1 Universal Validation Rules	27

		2.16.2 Validation Rules for Shader Capabilities	29
		2.16.3 Validation Rules for Kernel Capabilities	30
	2.17	Universal Limits	31
	2.18	Memory Model	31
		2.18.1 Memory Layout	32
		2.18.2 Aliasing	32
	2.19	Derivatives	32
	2.20	Code Motion	32
3	Rina	ry Form	33
	3.1	Magic Number	33
	3.2	Source Language	33
	3.3	Execution Model	33
	3.4		34
	3.5	Memory Model	34
	3.6	Execution Mode	34
	3.7	Storage Class	38
	3.8	Dim	39
	3.9	Sampler Addressing Mode	39
	3.10	Sampler Filter Mode	40
		Image Format	40
	3.12	Image Channel Order	41
	3.13	Image Channel Data Type	41
	3.14	Image Operands	42
	3.15	FP Fast Math Mode	44
	3.16	FP Rounding Mode	45
	3.17	Linkage Type	45
	3.18	Access Qualifier	45
	3.19	Function Parameter Attribute	46
	3.20	Decoration	46
	3.21	BuiltIn	51
	3.22	Selection Control	55
	3.23	Loop Control	55
	3.24	Function Control	56
	3.25	Memory Semantics <id></id>	56
	3.26	Memory Access	58
	3.27	Scope <id></id>	58
	3.28	Group Operation	59

	3.29	Kernel Enqueue Flags	60
	3.30	Kernel Profiling Info	61
	3.31	Capability	61
	3.32	Instructions	65
		3.32.1 Miscellaneous Instructions	65
		3.32.2 Debug Instructions	66
		3.32.3 Annotation Instructions	69
		3.32.4 Extension Instructions	71
		3.32.5 Mode-Setting Instructions	72
		3.32.6 Type-Declaration Instructions	74
		3.32.7 Constant-Creation Instructions	80
		3.32.8 Memory Instructions	85
		3.32.9 Function Instructions	89
		3.32.10 Image Instructions	91
		3.32.11 Conversion Instructions	11
		3.32.12 Composite Instructions	16
		3.32.13 Arithmetic Instructions	19
		3.32.14 Bit Instructions	27
		3.32.15 Relational and Logical Instructions	32
		3.32.16 Derivative Instructions	43
		3.32.17 Control-Flow Instructions	46
		3.32.18 Atomic Instructions	50
		3.32.19 Primitive Instructions	59
		3.32.20 Barrier Instructions	60
		3.32.21 Group Instructions	61
		3.32.22 Device-Side Enqueue Instructions	168
		3.32.23 Pipe Instructions	176
•	Char		105
A	Char		185
	A.1		185
	A.2	Changes from Version 1.00, Revision 1.	
	A.3	Changes from Version 1.00, Revision 1	
	A.4	Changes from Version 1.00, Revision 2	
	A.5	Changes from Version 1.00, Revision 3	.88

List of Tables

1	First Words of Physical Layout	18
2	Instruction Physical Layout	18
3	Limits	31

Contributors and Acknowledgments

Connor Abbott, Intel

Alexey Bader, Intel

Dan Baker, Oxide Games

Kenneth Benzie, Codeplay

Gordon Brown, Codeplay

Pat Brown, NVIDIA

Diana Po-Yu Chen, MediaTek

Stephen Clarke, Imagination

Patrick Doane, Blizzard Entertainment

Stefanus Du Toit, Google

Tim Foley, Intel

Ben Gaster, Qualcomm

Alexander Galazin, ARM

Christopher Gautier, ARM

Neil Henning, Codeplay

Kerch Holt, NVIDIA

Lee Howes, Qualcomm

Roy Ju, MediaTek

Daniel Koch, NVIDIA

Ashwin Kolhe, NVIDIA

Raun Krisch, Intel

Graeme Leese, Broadcom

Yuan Lin, NVIDIA

Yaxun Liu, AMD

Timothy Lottes, Epic Games

John McDonald, Valve

David Neto, Google

Christophe Riccio, Unity

Andrew Richards, Codeplay

Ian Romanick, Intel

Graham Sellers, AMD

Robert Simpson, Qualcomm

Brian Sumner, AMD

Andrew Woloszyn, Google

Weifeng Zhang, Qualcomm

Note

Up-to-date HTML and PDF versions of this specification may be found at the Khronos SPIR-V Registry. (https://www.khronos.org/registry/spir-v/)

1 Introduction

Abstract

SPIR-V is a simple binary intermediate language for graphical shaders and compute kernels. A SPIR-V module contains multiple entry points with potentially shared functions in the entry point's call trees. Each function contains a control-flow graph (CFG) of basic blocks, with optional instructions to express structured control flow. Load/store instructions are used to access declared variables, which includes all input/output (IO). Intermediate results bypassing load/store use static single-assignment (SSA) representation. Data objects are represented logically, with hierarchical type information: There is no flattening of aggregates or assignment to physical register banks, etc. Selectable addressing models establish whether general pointer operations may be used, or if memory access is purely logical.

This document fully defines **SPIR-V**, a Khronos-standard binary intermediate language for representing graphical-shader stages and compute kernels for multiple Khronos APIs.

1.1 Goals

SPIR-V has the following goals:

- Provide a simple binary intermediate language for all functionality appearing in Khronos shaders/kernels.
- Have a concise, transparent, self-contained specification (sections Specification and Binary Form).
- Map easily to other intermediate languages.
- Be the form passed by an API into a driver to set shaders/kernels.
- Can be targeted by new front ends for novel high-level languages.
- Allow the first steps of compilation and reflection to be done offline.
- Be low-level enough to require a reverse-engineering step to reconstruct source code.
- Improve portability by enabling shared tools to generate or operate on it.
- Allow separation of core specification from source-language-specific sets of built-in functions.
- Reduce compile time during application run time. (Eliminating most of the compile time during application run time is not a goal of this intermediate language. Target-specific register allocation and scheduling are still expected to take significant time.)
- Allow some optimizations to be done offline.

1.2 About this document

This document aims to:

- Include everything needed to fully understand, create, and consume SPIR-V. However:
 - Imported sets of instructions (which implement source-specific built-in functions) will need their own specification.
 - Many validation rules are client-API specific, and hence documented with client API and not in this specification.
- Separate expository and specification language. The specification-proper is in Specification and Binary Form.

1.3 Extendability

SPIR-V can be extended by multiple vendors or parties simultaneously:

- Using the OpExtension instruction to require new semantics that must be supported. Such new semantics would come from an extension document.
- Reserving (registering) ranges of the token values, as described further below.
- Aided by instruction skipping, also further described below.

Enumeration Token Values. It is easy to extend all the types, storage classes, opcodes, decorations, etc. by adding to the token values.

Registration. Ranges of token values in the Binary Form section can be pre-allocated to numerous vendors/parties. This allows combining multiple independent extensions without conflict. To register ranges, see https://www.khronos.org/registry/spir-v/api/spir-v.xml.

Extended Instructions. Sets of extended instructions can be provided and specified in separate specifications. These help personalize SPIR-V for different source languages or execution environments (client APIs). Multiple sets of extended instructions can be imported without conflict, as the extended instructions are selected by {set id, instruction number} pairs.

Instruction Skipping. Tools are encouraged to skip opcodes for features they are not required to process. This is trivially enabled by the word count in an instruction, which makes it easier to add new instructions without breaking existing tools.

1.4 Debuggability

SPIR-V can decorate, with a text string, virtually anything created in the shader: types, variables, functions, etc. This is required for externally visible symbols, and also allowed for naming the result of any instruction. This can be used to aid in understandability when disassembling or debugging lowered versions of SPIR-V.

Location information (file names, lines, and columns) can be interleaved with the instruction stream to track the origin of each instruction.

1.5 Design Principles

Regularity. All instructions start with a word count. This allows walking a SPIR-V module without decoding each opcode. All instructions have an opcode that dictates for all operands what kind of operand they are. For instructions with a variable number of operands, the number of variable operands is known by subtracting the number of non-variable words from the instruction's word count.

Non Combinatorial. There is no combinatorial type explosion or need for large encode/decode tables for types. Rather, types are parameterized. Image types declare their dimensionality, arrayness, etc. all orthogonally, which greatly simplify code. This is done similarly for other types. It also applies to opcodes. Operations are orthogonal to scalar/vector size, but not to integer vs. floating-point differences.

Modeless. After a given execution model (e.g., pipeline stage) is specified, internal operation is essentially modeless: Generally, it will follow the rule: "same spelling, same semantics", and does not have mode bits that modify semantics. If a change to SPIR-V modifies semantics, it should use a different spelling. This makes consumers of SPIR-V much more robust. There are execution modes declared, but these are generally to affect the way the module interacts with the environment around it, not the internal semantics. Capabilities are also declared, but this is to declare the subset of functionality that is used, not to change any semantics of what is used.

Declarative. SPIR-V declares externally-visible modes like "writes depth", rather than having rules that require deduction from full shader inspection. It also explicitly declares what addressing modes, execution model, extended instruction sets, etc. will be used. See Language Capabilities for more information.

SSA. All results of intermediate operations are strictly SSA. However, declared variables reside in memory and use load/store for access, and such variables can be stored to multiple times.

IO. Some storage classes are for input/output (IO) and, fundamentally, IO will be done through load/store of variables declared in these storage classes.

1.6 Static Single Assignment (SSA)

SPIR-V includes a phi instruction to allow the merging together of intermediate results from split control flow. This allows split control flow without load/store to memory. SPIR-V is flexible in the degree to which load/store is used; it is possible to use control flow with no phi-instructions, while still staying in SSA form, by using memory load/store.

Some storage classes are for IO and, fundamentally, IO will be done through load/store, and initial load and final store can never be eliminated. Other storage classes are shader local and can have their load/store eliminated. It can be considered an optimization to largely eliminate such loads/stores by moving them into intermediate results in SSA form.

1.7 Built-In Variables

SPIR-V identifies built-in variables from a high-level language with an enumerant decoration. This assigns any unusual semantics to the variable. Built-in variables must otherwise be declared with their correct SPIR-V type and treated the same as any other variable.

1.8 Specialization

Specialization enables creating a portable SPIR-V module outside the target execution environment, based on constant values that won't be known until inside the execution environment. For example, to size a fixed array with a constant not known during creation of a module, but known when the module will be lowered to the target architecture.

See Specialization in the next section for more details.

1.9 Example

The SPIR-V form is binary, not human readable, and fully described in Binary Form. This is an example disassembly to give a basic idea of what SPIR-V looks like:

GLSL fragment shader:

```
#version 450
in vec4 color1;
in vec4 multiplier;
noperspective in vec4 color2;
out vec4 color;
struct S {
   bool b;
   vec4 v[5];
   int i;
};
uniform blockName {
   S s;
   bool cond;
};
void main()
    vec4 scale = vec4(1.0, 1.0, 2.0, 1.0);
   if (cond)
        color = color1 + s.v[2];
    else
        color = sqrt(color2) * scale;
    for (int i = 0; i < 4; ++i)
        color *= multiplier;
```

Corresponding SPIR-V:

```
; Magic:
             0x07230203 (SPIR-V)
            0x00010000 (Version: 1.0.0)
; Version:
; Generator: 0x00080001 (Khronos Glslang Reference Front End; 1)
; Bound:
; Schema:
               OpCapability Shader
          %1 = OpExtInstImport "GLSL.std.450"
               OpMemoryModel Logical GLSL450
               OpEntryPoint Fragment %4 "main" %31 %33 %42 %57
               OpExecutionMode %4 OriginLowerLeft
; Debug information
               OpSource GLSL 450
               OpName %4 "main"
               OpName %9 "scale"
               OpName %17 "S"
               OpMemberName %17 0 "b"
               OpMemberName %17 1 "v"
               OpMemberName %17 2 "i"
```

```
OpName %18 "blockName"
              OpMemberName %18 0 "s"
              OpMemberName %18 1 "cond"
              OpName %20 ""
              OpName %31 "color"
              OpName %33 "color1"
              OpName %42 "color2"
              OpName %48 "i"
              OpName %57 "multiplier"
; Annotations (non-debug)
              OpDecorate %15 ArrayStride 16
              OpMemberDecorate %17 0 Offset 0
              OpMemberDecorate %17 1 Offset 16
              OpMemberDecorate %17 2 Offset 96
              OpMemberDecorate %18 0 Offset 0
              OpMemberDecorate %18 1 Offset 112
              OpDecorate %18 Block
              OpDecorate %20 DescriptorSet 0
              OpDecorate %42 NoPerspective
; All types, variables, and constants
         %2 = OpTypeVoid
                                                   ; void ()
         %3 = OpTypeFunction %2
                                                   ; 32-bit float
         %6 = OpTypeFloat 32
         %7 = OpTypeVector %6 4
                                                   ; vec4
         %8 = OpTypePointer Function %7 ; function-local vec4*
        %10 = OpConstant %6 1
        %11 = OpConstant %6 2
        %12 = OpConstantComposite %7 %10 %10 %11 %10; vec4(1.0, 1.0, 2.0, 1.0)
        %13 = OpTypeInt 32 0
                                                    ; 32-bit int, sign-less
        %14 = OpConstant %13 5
        %15 = OpTypeArray %7 %14
        %16 = OpTypeInt 32 1
        %17 = OpTypeStruct %13 %15 %16
        %18 = OpTypeStruct %17 %13
        %19 = OpTypePointer Uniform %18
        %20 = OpVariable %19 Uniform
        %21 = OpConstant %16 1
        %22 = OpTypePointer Uniform %13
        %25 = OpTypeBool
        %26 = OpConstant %13 0
        %30 = OpTypePointer Output %7
        %31 = OpVariable %30 Output
        %32 = OpTypePointer Input %7
        %33 = OpVariable %32 Input
        %35 = OpConstant %16 0
        %36 = OpConstant %16 2
        %37 = OpTypePointer Uniform %7
        %42 = OpVariable %32 Input
        %47 = OpTypePointer Function %16
        %55 = OpConstant %16 4
        %57 = OpVariable %32 Input
; All functions
         %4 = OpFunction %2 None %3
                                                     ; main()
         %5 = OpLabel
         %9 = OpVariable %8 Function
        %48 = OpVariable %47 Function
```

```
OpStore %9 %12
%23 = OpAccessChain %22 %20 %21 ; location of cond %24 = OpLoad %13 %23 ; load 32-bit int from cond %27 = OpINotEqual %25 %24 %26 ; convert to bool OpSelectionMerge %29 None ; structured if OpBranchConditional %27 %28 %41 ; if cond %28 = OpLabel ; then
%28 = OpLabel
                                                   ; then
%34 = OpLoad %7 %33
%38 = OpAccessChain %37 %20 %35 %21 %36 ; s.v[2]
%39 = OpLoad %7 %38
%40 = OpFAdd %7 %34 %39
      OpStore %31 %40
      OpBranch %29
%41 = OpLabel
                                                  ; else
%43 = OpLoad %7 %42
%44 = OpExtInst %7 %1 Sqrt %43 ; extended instruction sqrt
%45 = OpLoad %7 %9
%46 = OpFMul %7 %44 %45
      OpStore %31 %46
      OpBranch %29
%29 = OpLabel
                                                  ; endif
      OpStore %48 %35
      OpBranch %49
%49 = OpLabel
                                                 ; structured loop
      OpLoopMerge %51 %52 None
      OpBranch %53
%53 = OpLabel
%54 = OpLoad %16 %48
%56 = OpSLessThan %25 %54 %55 ; i < 4 ?
OpBranchConditional %56 %50 %51 ; body or break
%50 = OpLabel
                                                   ; body
%58 = OpLoad %7 %57
%59 = OpLoad %7 %31
%60 = OpFMul %7 %59 %58
      OpStore %31 %60
      OpBranch %52
                                                 ; continue target
%52 = OpLabel
%61 = OpLoad %16 %48
%62 = OpIAdd %16 %61 %21
                                       ; ++i
      OpStore %48 %62
      OpBranch %49
                                                   ; loop back
%51 = OpLabel
                                                   ; loop merge point
      OpReturn
       OpFunctionEnd
```

2 Specification

2.1 Language Capabilities

A SPIR-V module is consumed by an execution environment, specified by a client API, that needs to support the features used by that SPIR-V module. Features are classified through capabilities. Capabilities used by a particular SPIR-V module must be declared early in that module with the OpCapability instruction. Then:

- A validator can validate that the module uses only its declared capabilities.
- An execution environment is allowed to reject modules declaring capabilities it does not support. (See client API specifications for environment-specific rules.)

All available capabilities and their dependencies form a capability hierarchy, fully listed in the capability section. Only top-level capabilities need to be declared; their dependencies are automatically included.

This (SPIR-V) specification provides capability-specific validation rules, in the validation section. To ensure portability, each client API needs to include the following:

- Which capabilities in the capability section it requires environments to support, and hence allows in SPIR-V modules.
- Required limits, if they are beyond the Universal Limits.
- Any validation requirements specific to the environment that are not tied to specific capabilities, and hence not covered in the SPIR-V specification.

2.2 Terms

2.2.1 Instructions

Word: 32 bits.

< id >: A numerical name; the name used to refer to an object, a type, a function, a label, etc. An < id > always consumes one word. The < id > s defined by a module obey SSA.

Result <*id*>: Most instructions define a result, named by an <*id*> explicitly provided in the instruction. The *Result* <*id*> is used as an operand in other instructions to refer to the instruction that defined it.

Literal String: A nul-terminated stream of characters consuming an integral number of words. The character set is Unicode in the UTF-8 encoding scheme. The UTF-8 octets (8-bit bytes) are packed four per word, following the little-endian convention (i.e., the first octet is in the lowest-order 8 bits of the word). The final word contains the string's nul-termination character (0), and all contents past the end of the string in the final word are padded with 0.

Literal Number: A numeric value consuming one or more words. An instruction will determine what type a literal will be interpreted as. When the type's bit width is larger than one word, the literal's low-order words appear first. When the type's bit width is less than 32-bits, the literal's value appears in the low-order bits of the word, and the high-order bits must be 0 for a floating-point type, or 0 for an integer type with *Signedness* of 0, or sign extended when *Signedness* is 1. (Similarly for the remaining bits of widths larger than 32 bits but not a multiple of 32 bits.)

Literal: A Literal String or a Literal Number.

Operand: A one-word argument to an instruction. E.g., it could be an <id>, or a (part of a) literal. Which form it holds is always explicitly known from the opcode.

Immediate: Operand(s) directly holding a literal value rather than an <id>. Immediate values larger than one word will consume multiple operands, one per word. That is, operand counting is always done per word, not per immediate.

WordCount: The complete number of words taken by an instruction, including the word holding the word count and opcode, and any optional operands. An instruction's word count is the total space taken by the instruction.

Instruction: After a header, a module is simply a linear list of instructions. An instruction contains a word count, an opcode, an optional Result <id>, an optional <id> of the instruction's type, and a variable list of operands. All instruction opcodes and semantics are listed in Instructions.

Decoration: Auxiliary information such as built-in variable, stream numbers, invariance, interpolation type, relaxed precision, etc., added to <id>s or structure-type members through Decorations. Decorations are enumerated in Decoration in the Binary Form section.

Object: An instantiation of a non-void type, either as the Result <id> of an operation, or created through OpVariable.

Memory Object: An object created through OpVariable. Such an object can die on function exit, if it was a function variable, or exist for the duration of an entry point.

Intermediate Object or *Intermediate Value* or *Intermediate Result*: An object created by an operation (not memory allocated by OpVariable) and dying on its last consumption.

Constant Instruction: Either a specialization-constant instruction or a fixed constant instruction: Instructions that start "OpConstant" or "OpSpec".

[a, b]: This square-bracket notation means the range from a to b, inclusive of a and b. Parenthesis exclude their end point, so, for example, (a, b] means a to b excluding a but including b.

2.2.2 Types

Boolean type: The type returned by OpTypeBool.

Integer type: Any width signed or unsigned type from OpTypeInt. By convention, the lowest-order bit will be referred to as bit-number 0, and the highest-order bit as bit-number *Width* - 1.

Floating-point type: Any width type from OpTypeFloat.

Numerical type: An integer type or a floating-point type.

Scalar: A single instance of a numerical type or Boolean type. Scalars will also be called *components* when being discussed either by themselves or in the context of the contents of a vector.

Vector: An ordered homogeneous collection of two or more scalars. Vector sizes are quite restrictive and dependent on the execution model.

Matrix: An ordered homogeneous collection of vectors. When vectors are part of a matrix, they will also be called *columns*. Matrix sizes are quite restrictive and dependent on the execution model.

Array: An ordered homogeneous collection of any non-void-type objects. When an object is part of an array, it will also be called an *element*. Array sizes are generally not restricted.

Structure: An ordered heterogeneous collection of any non-void types. When an object is part of a structure, it will also be called a *member*.

Aggregate: A structure or an array.

Composite: An aggregate, a matrix, or a vector.

Image: A traditional texture or image; SPIR-V has this single name for these. An image type is declared with OpTypeImage. An image does not include any information about how to access, filter, or sample it.

Sampler: Settings that describe how to access, filter, or sample an image. Can come either from literal declarations of settings or be an opaque reference to externally bound settings. A sampler does not include an image.

Sampled Image: An image combined with a sampler, enabling filtered accesses of the image's contents.

Concrete Type: A numerical scalar, vector, or matrix type, or OpTypePointer when using a **Physical** addressing model, or any aggregate containing only these types.

Abstract Type: An OpTypeVoid or OpTypeBool, or OpTypePointer when using the **Logical** addressing model, or any aggregate type containing any of these.

Opaque Type: A type that is, or contains, or points to, or contains pointers to, any of the following types:

- OpTypeImage
- OpTypeSampler
- OpTypeSampledImage
- OpTypeOpaque
- OpTypeEvent
- OpTypeDeviceEvent
- OpTypeReserveId
- OpTypeQueue
- OpTypePipe
- OpTypeForwardPointer

2.2.3 Module

Module: A single unit of SPIR-V. It can contain multiple entry points, but only one set of capabilities.

Entry Point: A function in a module where execution begins. A single *entry point* is limited to a single execution model. An entry point is declared using OpEntryPoint.

Execution Model: A graphical-pipeline stage or OpenCL kernel. These are enumerated in Execution Model.

Execution Mode: Modes of operation relating to the interface or execution environment of the module. These are enumerated in Execution Mode. Generally, modes do not change the semantics of instructions within a SPIR-V module.

Vertex Processor: Any stage or execution model that processes vertices: Vertex, tessellation control, tessellation evaluation, and geometry. Explicitly excludes fragment and compute execution models.

2.2.4 Control Flow

Block: A contiguous sequence of instructions starting with an OpLabel, ending with a branch instruction, and having no other label or branch instructions.

Branch Instruction: One of the following, used to terminate blocks:

- OpBranch
- OpBranchConditional
- OpSwitch
- OpKill
- OpReturn
- OpReturnValue
- OpUnreachable

Dominate: A block A dominates a block B, where A and B are in the same function, if every path from the function's entry point to block B goes through block A.

Post Dominate: A block *B* post dominates a block *A*, where *A* and *B* are in the same function, if every path from *A* to a function-return instruction goes through block *B*.

Control-Flow Graph: The graph formed by a function's blocks and branches. The blocks are the graph's nodes, and the branches the graph's edges.

CFG: Control-flow graph.

Back Edge: If a depth-first traversal is done on a function's CFG, starting from the first block of the function, a back edge is a branch to a previously visited block. A back-edge block is the block containing such a branch.

Merge Instruction: One of the following, used before a branch instruction to declare structured control flow:

- OpSelectionMerge
- OpLoopMerge

Header Block: A block containing a merge instruction.

Loop Header: A header block whose merge instruction is an OpLoopMerge.

Merge Block: A block declared by the Merge Block operand of a merge instruction.

Break Block: A block containing a branch to the Merge Block of a loop header's merge instruction.

Continue Block: A block containing a branch to an OpLoopMerge instruction's Continue Target.

Return Block: A block containing an OpReturn or OpReturn Value branch.

Invocation: A single execution of an entry point in a SPIR-V module, operating only on the amount of data explicitly exposed by the semantics of the instructions. (Any implicit operation on additional instances of data would comprise additional invocations.) For example, in compute execution models, a single invocation operates only on a single work item, or, in a vertex execution model, a single invocation operates only on a single vertex.

Subgroup: The set of invocations exposed as running concurrently with the current invocation. In compute models, the current workgroup is a superset of the subgroup.

Invocation Group: The complete set of invocations collectively processing a particular compute workgroup or graphical operation, where the scope of a "graphical operation" is implementation dependent, but at least as large as a single point, line, triangle, or patch, and at most as large as a single rendering command, as defined by the client API.

Derivative Group: Defined only for the **Fragment** Execution Model: The set of invocations collectively processing a single point, line, or triangle, including any helper invocations.

Dynamic Instance: Within a single invocation, a single static instruction can be executed multiple times, giving multiple dynamic instances of that instruction. This can happen when the instruction is executed in a loop, or in a function called from multiple call sites, or combinations of multiple of these. Different loop iterations and different dynamic function-call-site chains yield different dynamic instances of such an instruction. Dynamic instances are distinguished by the control-flow path within an invocation, not by which invocation executed it. That is, different invocations of an entry point execute the same dynamic instances of an instruction when they follow the same control-flow path, starting from that entry point.

Dynamically Uniform: An <id> is dynamically uniform for a dynamic instance consuming it when its value is the same for all invocations (in the invocation group) that execute that dynamic instance.

Uniform Control Flow: Uniform control flow (or converged control flow) occurs when all invocations in the invocation group or derivative group execute the same control-flow path (and hence the same sequence of dynamic instances of instructions). Uniform control flow is the initial state at the entry point, and lasts until a conditional branch takes different control paths for different invocations (non-uniform or divergent control flow). Such divergence can reconverge, with all the invocations once again executing the same control-flow path, and this re-establishes the existence of uniform control flow. If control flow is uniform upon entry into a header block, and all invocations leave that dynamic instance of the header block's control-flow construct via the header block's declared merge block, then control flow reconverges to be uniform at that merge block.

2.3 Physical Layout of a SPIR-V Module and Instruction

A SPIR-V module is a single linear stream of words. The first words are shown in the following table:

Table 1: First Words of Physical Layout

Word	Contents	
Number		
0	Magic Number.	
1	Version number. The bytes are, high-order to low-order:	
	0 Major Number Minor Number 0	
	Hence, version 1.00 is the value 0x00010000.	
2	Generator's magic number. It is associated with the tool that generated	
	the module. Its value does not affect any semantics, and is allowed to be	
	0. Using a non-0 value is encouraged, and can be registered with	
	Khronos at https://www.khronos.org/registry/spir-v/api/spir-v.xml.	
3	Bound; where all <id>s in this module are guaranteed to satisfy</id>	
0 < id < Bound		
	Bound should be small, smaller is better, with all <id> in a module being</id>	
	densely packed and near 0.	
4	0 (Reserved for instruction schema, if needed.)	
5	First word of instruction stream, see below.	

All remaining words are a linear sequence of instructions.

Each instruction is a stream of words:

Table 2: Instruction Physical Layout

Instruction	Contents	
Word Number		
0	Opcode: The 16 high-order bits are the WordCount of the	
	instruction. The 16 low-order bits are the opcode enumerant.	
1	Optional instruction type <id> (presence determined by opcode).</id>	
	Optional instruction Result <id> (presence determined by</id>	
	opcode).	
	Operand 1 (if needed)	
	Operand 2 (if needed)	
• • •		
WordCount - 1 Operand N (N is determined by WordCount minus the 1 to 3		
	words used for the opcode, instruction type <i><id></id></i> , and instruction	
	Result $\langle id \rangle$).	

Instructions are variable length due both to having optional instruction type <*id*> and *Result* <*id*> words as well as a variable number of operands. The details for each specific instruction are given in the Binary Form section.

2.4 Logical Layout of a Module

The instructions of a SPIR-V module must be in the following order:

- 1. All OpCapability instructions.
- 2. Optional OpExtension instructions (extensions to SPIR-V).
- 3. Optional OpExtInstImport instructions.
- 4. The single required OpMemoryModel instruction.
- 5. All entry point declarations, using OpEntryPoint.
- 6. All execution mode declarations, using OpExecutionMode.
- 7. These debug instructions, which must be in the following order:
 - a. all OpString, OpSourceExtension, OpSource, and OpSourceContinued, without forward references.
 - b. all OpName and all OpMemberName
- 8. All annotation instructions:
 - a. all decoration instructions (OpDecorate, OpMemberDecorate, OpGroupDecorate, OpGroupMemberDecorate, and OpDecorationGroup).
- 9. All type declarations (OpTypeXXX instructions), all constant instructions, and all global variable declarations (all OpVariable instructions whose Storage Class is not Function). All operands in all these instructions must be declared before being used. Otherwise, they can be in any order. This section is also the first section to allow use of OpLine debug information.
- 10. All function declarations ("declarations" are functions without a body; there is no forward declaration to a function with a body). A function declaration is as follows.
 - a. Function declaration, using OpFunction.
 - b. Function parameter declarations, using OpFunctionParameter.
 - c. Function end, using OpFunctionEnd.
- 11. All function definitions (functions with a body). A function definition is as follows.
 - a. Function definition, using OpFunction.
 - b. Function parameter declarations, using OpFunctionParameter.
 - c. Block
 - d. Block
 - e. ...
 - f. Function end, using OpFunctionEnd.

Within a function definition:

- A block always starts with an OpLabel instruction. This may be immediately preceded by an OpLine instruction, but the **OpLabel** is considered as the beginning of the block.
- A block always ends with a branch instruction (see validation rules for more detail).
- All OpVariable instructions in a function must have a Storage Class of Function.
- All OpVariable instructions in a function must be in the first block in the function. These instructions, together with any immediately preceding OpLine instructions, must be the first instructions in that block. (Note the validation rules prevent OpPhi instructions in the first block of a function.)
- A function definition (starts with OpFunction) can be immediately preceded by an OpLine instruction.

Forward references (an operand <id> that appears before the Result <id> defining it) are allowed for:

- Operands that are an OpFunction. This allows for recursion and early declaration of entry points.
- Annotation-instruction operands. This is required to fully know everything about a type or variable once it is declared.
- · Labels.
- Loops can have forward references to a phi function.
- An OpTypeForwardPointer has a forward reference to an OpTypePointer.
- An OpTypeStruct operand that's a forward reference to the *Pointer Type* operand to an OpTypeForwardPointer.
- The list of *<id>* provided in the OpEntryPoint instruction.

In all cases, there is enough type information to enable a single simple pass through a module to transform it. For example, function calls have all the type information in the call, phi-functions don't change type, and labels don't have type. The pointer forward reference allows structures to contain pointers to themselves or to be mutually recursive (through pointers), without needing additional type information.

The Validation Rules section lists additional rules that must be satisfied.

2.5 Instructions

Most instructions create a Result <id>, as provided in the Result <id> field of the instruction. These Result <id>s are then referred to by other instructions through their <id> operands. All instruction operands are specified in the Binary Form section.

Instructions are explicit about whether they require immediates, rather than an $\langle id \rangle$ referring to some other result. This is strictly known just from the opcode.

- An immediate 32-bit (or smaller) integer is always one operand directly holding a 32-bit two's-complement value.
- An immediate 32-bit float is always one operand, directly holding a 32-bit IEEE 754 floating-point representation.
- An immediate 64-bit float is always two operands, directly holding a 64-bit IEEE 754 representation. The low-order 32 bits appear in the first operand.

2.5.1 SSA Form

A module is always in static single assignment (SSA) form. That is, there is always exactly one instruction resulting in any particular Result <id>. Storing into variables declared in memory is not subject to this; such stores do not create Result <id>s. Accessing declared variables is done through:

- OpVariable to allocate an object in memory and create a Result <id> that is the name of a pointer to it.
- OpAccessChain or OpInBoundsAccessChain to create a pointer to a subpart of a composite object in memory.
- OpLoad through a pointer, giving the loaded object a Result <id> that can then be used as an operand in other instructions.
- OpStore through a pointer, to write a value. There is no Result <id> for an OpStore.

OpLoad and OpStore instructions can often be eliminated, using intermediate results instead. When this happens in multiple control-flow paths, these values need to be merged again at the path's merge point. Use OpPhi to merge such values together.

2.6 Entry Point and Execution Model

The OpEntryPoint instruction identifies an entry point with two key things: an execution model and a function definition. Execution models include **Vertex**, **GLCompute**, etc. (one for each graphical stage), as well as **Kernel** for OpenCL kernels. For the complete list, see Execution Model. An OpEntryPoint also supplies a name that can be used externally to identify the entry point, and a declaration of all the **Input** and **Output** variables that form its input/output interface.

The static function call graphs rooted at two entry points are allowed to overlap, so that function definitions and global variable definitions can be shared. The execution model and any execution modes associated with an entry point apply to the entire static function call graph rooted at that entry point. This rule implies that a function appearing in both call graphs of two distinct entry points may behave differently in each case. Similarly, variables whose semantics depend on properties of an entry point, e.g. those using the **Input Storage Class**, may behave differently when used in call graphs rooted in two different entry points.

2.7 Execution Modes

Information like the following is declared with OpExecutionMode instructions. For example,

- number of invocations (Invocations)
- vertex-order CCW (VertexOrderCcw)
- triangle strip generation (OutputTriangleStrip)
- number of output vertices (OutputVertices)
- etc.

For a complete list, see Execution Mode.

2.8 Types and Variables

Types are built up hierarchically, using OpTypeXXX instructions. The Result <id> of an OpTypeXXX instruction becomes a type <id> for future use where type <id>s are needed (therefore, OpTypeXXX instructions do not have a type <id>, like most other instructions do).

The "leaves" to start building with are types like OpTypeFloat, OpTypeInt, OpTypeImage, OpTypeEvent, etc. Other types are built up from the *Result <id>* of these. The numerical types are parameterized to specify bit width and signed vs. unsigned.

Higher-level types are then constructed using opcodes like OpTypeVector, OpTypeMatrix, OpTypeImage, OpTypeArray, OpTypeRuntimeArray, OpTypeStruct, and OpTypePointer. These are parameterized by number of components, array size, member lists, etc. The image types are parameterized by the return type, dimensionality, arrayness, etc. To do sampling or filtering operations, a type from OpTypeSampledImage is used that contains both an image and a sampler. Such a sampled image can be set directly by the API, or combined in a SPIR-V module from an independent image and an independent sampler.

Types are built bottom up: A parameterizing operand in a type must be defined before being used.

Some additional information about the type of an <id> can be provided using the decoration instructions (OpDecorate, OpMemberDecorate, OpGroupMemberDecorate, and OpDecorationGroup). These can add, for example, **Invariant** to an <id> created by another instruction. See the full list of Decorations in the Binary Form section.

Two different type < id>s form, by definition, two different types. It is valid to declare multiple aggregate type < id>s having the same opcode and operands. This is to allow multiple instances of aggregate types with the same structure to be decorated differently. (Different decorations are not required; two different aggregate type < id>s are allowed to have identical declarations and decorations, and will still be two different types.) Non-aggregate types are different: It is invalid to declare multiple type < id>s for the same scalar, vector, or matrix type. That is, non-aggregate type declarations must all have different opcodes or operands. (Note that non-aggregate types cannot be decorated in ways that affect their type.)

Variables are declared to be of an already built type, and placed in a Storage Class. Storage classes include **UniformConstant**, **Input**, **Workgroup**, etc. and are fully specified in Storage Class. Variables declared with the **Function** Storage Class can have their lifetime's specified within their function using the OpLifetimeStart and OpLifetimeStop instructions.

Intermediate results are typed by the instruction's type <id>, which must validate with respect to the operation being done.

Built-in variables needing special driver handling (having unique semantics) are declared using OpDecorate or OpMemberDecorate with the **BuiltIn** Decoration, followed by a BuiltIn enumerant. This decoration is applied to a variable or a structure-type member.

2.9 Function Calling

To call a function defined in the current module or a function declared to be imported from another module, use OpFunctionCall with an operand that is the <id> of the OpFunction to call, and the <id>s of the arguments to pass. All arguments are passed by value into the called function. This includes pointers, through which a callee object could be modified.

2.10 Extended Instruction Sets

Many operations and/or built-in function calls from high-level languages are represented through *extended instruction sets*. Extended instruction sets will include things like

- trigonometric functions: sin(), cos(), ...
- exponentiation functions: exp(), pow(), ...
- geometry functions: reflect(), smoothstep(), ...
- functions having rich performance/accuracy trade-offs
- etc

Non-extended instructions, those that are core SPIR-V instructions, are listed in the Binary Form section. Native operations include:

- Basic arithmetic: +, -, *, min(), scalar * vector, etc.
- Texturing, to help with back-end decoding and support special code-motion rules.
- Derivatives, due to special code-motion rules.

Extended instruction sets are specified in independent specifications. They can be referenced (but not specified) in this specification. The separate extended instruction set specification will specify instruction opcodes, semantics, and instruction names.

To use an extended instruction set, first import it by name string using OpExtInstImport and giving it a Result <id>:

```
<extinst-id> OpExtInstImport "name-of-extended-instruction-set"
```

The "name-of-extended-instruction-set" is a literal string. The standard convention for this string is

```
"<source language name>.<package name>.<version>"
```

For example "GLSL.std.450" could be the name of the core built-in functions for GLSL versions 450 and earlier.

Note

There is nothing precluding having two "mirror" sets of instructions with different names but the same opcode values, which could, for example, let modifying just the import statement to change a performance/accuracy trade off.

Then, to call a specific extended instruction, use OpExtInst:

```
OpExtInst <extinst-id> instruction-number operand0, operand1, ...
```

Extended instruction-set specifications will provide semantics for each "instruction-number". It is up to the specific specification what the overloading rules are on operand type. The specification must be clear on its semantics, and producers/consumers of it must follow those semantics.

By convention, it is recommended that all external specifications include an **enum** {...} listing all the "instruction-numbers", and a mapping between these numbers and a string representing the instruction name. However, there are no requirements that instruction name strings are provided or mangled.

Note

Producing and consuming extended instructions can be done entirely through numbers (no string parsing). An extended instruction set specification provides opcode enumerant values for the instructions, and these will be produced by the front end and consumed by the back end.

2.11 Structured Control Flow

SPIR-V can explicitly declare structured control-flow *constructs* using merge instructions. These explicitly declare a header block before the control flow diverges and a merge block where control flow subsequently converges. These blocks delimit constructs that must nest, and can only be entered and exited in structured ways, as per the following.

Structured control-flow declarations must satisfy the following rules:

- the merge block declared by a header block cannot be a merge block declared by any other header block
- each header block must dominate its merge block, unless the merge block is unreachable in the CFG
- all CFG back edges must branch to a loop header, with each loop header having exactly one back edge branching to it
- for a given loop, its back-edge block must post dominate the OpLoopMerge's *Continue Target*, and that *Continue Target* must dominate that back-edge block

A structured control-flow construct is then defined as one of:

- a *selection construct*: the set of blocks dominated by a selection header, minus the set of blocks dominated by the header's merge block
- a *continue construct*: the set of blocks dominated by an OpLoopMerge's *Continue Target* and post dominated by the corresponding back-edge block
- a *loop construct*: the set of blocks dominated by a *loop header*, minus the set of blocks dominated by the loop's merge block, minus the loop's corresponding *continue construct*
- a case construct: the set of blocks dominated by an OpSwitch Target or Default, minus the set of blocks dominated by the OpSwitch's merge block (this construct is only defined for those OpSwitch Target or Default that are not equal to the OpSwitch's corresponding merge block)

The above structured control-flow constructs must satisfy the following rules:

- if a construct contains another header block, then it also contains that header's corresponding merge block
- the only blocks in a construct that can branch outside the construct are
 - a block branching to the construct's merge block
 - a block branching from one case construct to another, for the same **OpSwitch**
 - a continue block for the innermost loop it is nested inside of

- a break block for the innermost loop it is nested inside of
- a return block
- additionally for switches:
 - an **OpSwitch** block dominates all its defined *case constructs*
 - each case construct has at most one branch to another case construct
 - each *case construct* is branched to by at most one other *case construct*
 - if Target T1 branches to Target T2, or if Target T1 branches to the Default and the Default branches to Target T2, then
 T1 must immediately precede T2 in the list of the OpSwitch Target operands

2.12 Specialization

Specialization is intended for constant objects that will not have known constant values until after initial generation of a SPIR-V module. Such objects are called *specialization constants*.

A SPIR-V module containing specialization constants can consume one or more externally provided *specializations*: A set of final constant values for some subset of the module's *specialization constants*. Applying these final constant values yields a new module having fewer remaining specialization constants. A module also contains default values for any specialization constants that never get externally specialized.

Note

No optimizing transforms are required to make a *specialized* module functionally correct. The specializing transform is straightforward and explicitly defined below.

Note

Ad hoc specializing should not be done through constants (OpConstant or OpConstantComposite) that get overwritten: A SPIR-V \rightarrow SPIR-V transform might want to do something irreversible with the value of such a constant, unconstrained from the possibility that its value could be later changed.

Within a module, a *Specialization Constant* is declared with one of these instructions:

- OpSpecConstantTrue
- OpSpecConstantFalse
- OpSpecConstant
- OpSpecConstantComposite
- OpSpecConstantOp

The literal operands to OpSpecConstant are the default numerical specialization constants. Similarly, the "True" and "False" parts of OpSpecConstantTrue and OpSpecConstantFalse provide the default Boolean specialization constants. These default values make an external specialization optional. However, such a default constant is applied only after all external specializations are complete, and none contained a specialization for it.

An external specialization is provided as a logical list of pairs. Each pair is a **SpecId** Decoration of a scalar specialization instruction along with its specialization constant. The numeric values are exactly what the operands would be to a corresponding OpConstant instruction. Boolean values are true if non-zero and false if zero.

Specializing a module is straightforward. The following specialization-constant instructions can be updated with specialization constants, and replaced in place, leaving everything else in the module exactly the same:

```
OpSpecConstantTrue -> OpConstantTrue or OpConstantFalse
   OpSpecConstantFalse -> OpConstantTrue or OpConstantFalse
        OpSpecConstant -> OpConstant
OpSpecConstantComposite -> OpConstantComposite
```

The OpSpecConstantOp instruction is specialized by executing the operation and replacing the instruction with the result. The result can be expressed in terms of a constant instruction that is not a specialization-constant instruction. (Note, however, this resulting instruction might not have the same size as the original instruction, so is not a "replaced in place" operation.)

When applying an external specialization, the following (and only the following) must be modified to be non-specialization-constant instructions:

- specialization-constant instructions with values provided by the specialization
- specialization-constant instructions that consume nothing but non-specialization constant instructions (including those
 that the partial specialization transformed from specialization-constant instructions; these are in order, so it is a single
 pass to do so)

A full specialization can also be done, when requested or required, in which all specialization-constant instructions will be modified to non-specialization-constant instructions, using the default values where required.

2.13 Linkage

The ability to have partially linked modules and libraries is provided as part of the Linkage capability.

By default, functions and global variables are private to a module and cannot be accessed by other modules. However, a module may be written to *export* or *import* functions and global (module scope) variables. Imported functions and global variable definitions are resolved at linkage time. A module is considered to be partially linked if it depends on imported values.

Within a module, imported or exported values are decorated using the **Linkage Attributes** Decoration. This decoration assigns the following linkage attributes to decorated values:

- A Linkage Type.
- A name, which is a Literal String, and is used to uniquely identify exported values.

Note

When resolving imported functions, the Function Control and all Function Parameter Attributes are taken from the function definition, and not from the function declaration.

2.14 Relaxed Precision

The **RelaxedPrecision** Decoration allows 32-bit integer and 32-bit floating-point operations to execute with a relaxed precision of somewhere between 16 and 32 bits.

For a floating-point operation, operating at relaxed precision means that the minimum requirements for range and precision are as follows:

- the floating point range may be as small as $(-2^{14}, 2^{14})$
- the floating point magnitude range may be as small as $(2^{-14}, 2^{14})$
- the relative floating point precision may be as small as 2⁻¹⁰

Relative floating-point precision is defined as the worst case (i.e. largest) ratio of the smallest step in relation to the value for all non-zero values:

Precision_{relative} = $(abs(v_1 - v_2)_{min} / abs(v_1))_{max}$ for $v_1 \neq 0$, $v_2 \neq 0$, $v_1 \neq v_2$

For integer operations, operating at relaxed precision means that the operation will be evaluated by an operation in which, for some N, $16 \le N \le 32$:

- all inputs are truncated to N bits, using either signed or unsigned truncation as appropriate for the operation in question,
- the operation is executed as though its type were N bits in size,
- finally, the result is zero or sign extended to 32 bits as determined by the signedness of the result type of the operation.

The **RelaxedPrecision** Decoration can be applied to:

- The <id> of a variable, where the variable's type is a scalar, vector, or matrix, or an array of scalar, vector, or matrix. In all cases, the components in the type must be a 32-bit numerical type.
- The Result <id> of an instruction that operates on numerical types, meaning the instruction is to operate at relaxed precision.
- The Result <id> of an OpFunction meaning the function's returned result is at relaxed precision. It cannot be applied to OpTypeFunction or to an **OpFunction** whose return type is **OpTypeVoid**.
- A structure-type member (through OpMemberDecorate).

When applied to a variable or structure member, all loads and stores from the decorated object may be treated as though they were decorated with **RelaxedPrecision**. Loads may also be decorated with **RelaxedPrecision**, in which case they are treated as operating at relaxed precision.

All loads and stores involving relaxed precision still read and write 32 bits of data, respectively. Floating-point data read or written in such a manner is written in full 32-bit floating-point format. However, a load or store might reduce the precision (as allowed by **RelaxedPrecision**) of the destination value.

For debugging portability of floating-point operations, OpQuantizeToF16 may be used to explicitly reduce the precision of a relaxed-precision result to 16-bit precision. (Integer-result precision can be reduced, for example, using left- and right-shift opcodes.)

2.15 Debug Information

Debug information is supplied with:

- Source-code text through OpString, OpSource, and OpSourceContinued.
- Object names through OpName and OpMemberName.
- Line numbers through OpLine.

A module will not lose any semantics when all such instructions are removed.

2.15.1 Function-Name Mangling

There is no functional dependency on how functions are named. Signature-typing information is explicitly provided, without any need for name "unmangling". (Valid modules can be created without inclusion of mangled names.)

By convention, for debugging purposes, modules with OpSource Source Language of OpenCL use the Itanium name-mangling standard.

2.16 Validation Rules

2.16.1 Universal Validation Rules

All modules must obey the following, or it is an invalid module:

- The stream of instructions must be ordered as described in the Logical Layout section.
- Any use of a feature described by a capability in the capability section requires that capability to be declared, either directly, or as a "depends on" capability on a capability that is declared.
- Non-structure types (scalars, vectors, arrays, etc.) with the same operand parameterization cannot be type aliases. For non-structures, two type <*id*>*s* match if-and-only-if the types match.
- If the **Logical** addressing model is selected:
 - OpVariable cannot allocate an object whose type is a pointer type (that is, it cannot create an object in memory that is
 itself a pointer and whose result would thus be a pointer to a pointer)
 - A pointer can only be an operand to the following instructions
 - * OpLoad
 - * OpStore
 - * OpAccessChain
 - * OpInBoundsAccessChain
 - A pointer can only be created by the following instructions:
 - * OpVariable
 - * OpAccessChain
 - * OpInBoundsAccessChain
 - All indexes in OpAccessChain and OpInBoundsAccessChain that are OpConstant with type of OpTypeInt with a signedness of 1 must not have their sign bit set.

• SSA

- Each <id> must appear exactly once as the Result <id> of an instruction.
- The definition of an SSA <id> should dominate all uses of it, with the following exceptions:
 - * Function calls may call functions not yet defined. However, note that the function's argument and return types will already be known at the call site.
 - * Uses in a phi-function in a loop may consume definitions in the loop that don't dominate the use.
- Entry point and execution model
 - There is at least one OpEntryPoint instruction, unless the Linkage capability is being used.
 - No function can be targeted by both an OpEntryPoint instruction and an OpFunctionCall instruction.
- Functions
 - A function declaration (an OpFunction with no basic blocks), must have a Linkage Attributes Decoration with the Import Linkage Type.
 - A function definition (an OpFunction with basic blocks) cannot be decorated with the **Import** Linkage Type.
 - A function cannot have both a declaration and a definition (no forward declarations).
- Global (Module Scope) Variables
 - It is illegal to initialize an imported variable. This means that a module-scope OpVariable with initialization value cannot be marked with the **Import** Linkage Type.
- Control-Flow Graph (CFG)

- Blocks exist only within a function.
- The first block in a function definition is the entry point of that function and cannot be the target of any branch. (Note this means it will have no OpPhi instructions.)
- The order of blocks in a function must satisfy the rule that blocks appear before all blocks they dominate.
- Each block starts with a label.
 - * A label is made by OpLabel.
 - * This includes the first block of a function (**OpFunction** is not a label).
 - * Labels are used only to form blocks.
- The last instruction of each block is a branch instruction.
- Branch instructions can only appear as the last instruction in a block.
- OpLabel instructions can only appear within a function.
- All branches within a function must be to labels in that function.
- All OpFunctionCall Function operands are an <id> of an OpFunction in the same module.
- · Data rules
 - Scalar floating-point types can be parameterized only as 32 bit, plus any additional sizes enabled by capabilities.
 - Scalar integer types can be parameterized only as 32 bit, plus any additional sizes enabled by capabilities.
 - Vector types can only be parameterized with numerical types or the OpTypeBool type.
 - Vector types for can only be parameterized as having 2, 3, or 4 components, plus any additional sizes enabled by capabilities.
 - Matrix types can only be parameterized with floating-point types.
 - Matrix types can only be parameterized as having only 2, 3, or 4 columns.
 - Specialization constants (see Specialization) are limited to integers, Booleans, floating-point numbers, and vectors of these.
 - Forward reference operands in an OpTypeStruct
 - * must be later declared with OpTypePointer
 - * the type pointed to must be an OpTypeStruct
 - * had an earlier OpTypeForwardPointer forward reference to the same <id>
 - All OpSampledImage instructions must be in the same block in which their Result <id> are consumed. Result <id> from OpSampledImage instructions must not appear as operands to OpPhi instructions or OpSelect instructions, or any instructions other than the image lookup and query instructions specified to take an operand whose type is OpTypeSampledImage.
 - Instructions for extracting a scalar image or scalar sampler out of a composite must only use dynamically-uniform indexes. They must be in the same block in which their *Result <id>* are consumed. Such *Result <id>* must not appear as operands to OpPhi instructions or OpSelect instructions, or any instructions other than the image instructions specified to operate on them.
- · Decoration rules
 - The **Aliased** Decoration can only be applied to intermediate objects that are pointers to non-void types.
 - The Linkage Attributes Decoration cannot be applied to functions targeted by an OpEntryPoint instruction.
 - A BuiltIn Decoration can only be applied as follows:
 - * When applied to a structure-type member, all members of that structure type must also be decorated with **BuiltIn**. (No allowed mixing of built-in variables and non-built-in variables within a single structure.)
 - * When applied to a structure-type member, that structure type cannot be contained as a member of another structure type.
 - * There is at most one object per Storage Class that can contain a structure type containing members decorated with **BuiltIn**, consumed per entry-point.
- OpLoad and OpStore can only consume objects whose type is a pointer.

- A Result <id> resulting from an instruction within a function can only be used in that function.
- A function call must have the same number of arguments as the function definition (or declaration) has parameters, and their respective types must match.
- · An instruction requiring a specific number of operands must have that many operands. The word count must agree.
- Each opcode specifies its own requirements for number and type of operands, and these must be followed.
- Atomic access rules
 - The pointers taken by atomic operation instructions must be a pointer into one of the following Storage Classes:
 - * Uniform when used with the BufferBlock Decoration
 - * Workgroup
 - * CrossWorkgroup
 - * Generic
 - * AtomicCounter
 - * Image
 - The only instructions that can operate on a pointer to the **AtomicCounter Storage Class** are
 - * OpAtomicLoad
 - * OpAtomicIIncrement
 - * OpAtomicIDecrement
 - All pointers used in atomic operation instructions must be pointers to one of the following:
 - * 32-bit scalar integer
 - * 64-bit scalar integer

2.16.2 Validation Rules for Shader Capabilities

- · CFG:
 - Loops must be structured, having an OpLoopMerge instruction in their header.
 - Selections must be structured, having an OpSelectionMerge instruction in their header.
- · Entry point and execution model
 - Each entry point in a module, along with its corresponding static call tree within that module, forms a complete pipeline stage.
 - Each OpEntryPoint with the Fragment Execution Model must have an OpExecutionMode for either the OriginLowerLeft or the OriginUpperLeft Execution Mode. (Exactly one of these is required.)
 - An OpEntryPoint with the Fragment Execution Model can set at most one of the DepthGreater, DepthLess, or DepthUnchanged Execution Modes.
 - An OpEntryPoint with one of the Tessellation Execution Modes can set at most one of the SpacingEqual, FractionalEven, or FractionalOdd Execution Modes.
 - An OpEntryPoint with one of the Tessellation Execution Models can set at most one of the Triangles, Quads, or Isolines Execution Modes.
 - An OpEntryPoint with one of the Tessellation Execution Models can set at most one of the VertexOrderCw or VertexOrderCcw Execution Modes.
 - An OpEntryPoint with the Geometry Execution Model must set exactly one of the InputPoints, InputLines, InputLinesAdjacency, Triangles, or TrianglesAdjacency Execution Modes.
 - An OpEntryPoint with the Geometry Execution Model must set exactly one of the OutputPoints, OutputLineStrip, or OutputTriangleStrip Execution Modes.
- Composite objects in the **UniformConstant**, **Uniform**, and **PushConstant** Storage Classes must be explicitly laid out. The following apply to all the aggregate and matrix types describing such an object, recursively through their nested types:

- Each structure-type member must have an **Offset** Decoration.
- Each array type must have an **ArrayStride** Decoration.
- Each structure-type member that is a matrix or array-of-matrices must have be decorated with
 - * a MatrixStride Decoration, and
 - * one of the **RowMajor** or **ColMajor** Decorations.
- The **ArrayStride**, **MatrixStride**, and **Offset** Decorations must be large enough to hold the size of the objects they affect (that is, specifying overlap is invalid).
- The **MatrixStride** on a **RowMajor** (**ColMajor**) matrix must be padded to hold a row (column) of 4 components, when the matrix only has 3 columns (rows). In all other uses of **MatrixStride**, no padding is allowed.
- For structure objects in the **Input** and **Output** Storage Classes, the following apply:
 - When applied to structure-type members, the Decorations Noperspective, Flat, Patch, Centroid, and Sample can
 only be applied to the top-level members of the structure type. (Nested objects' types cannot be structures whose
 members are decorated with these decorations.)
- · Decorations
 - At most one of **Noperspective** or **Flat** Decorations can be applied to the same object or member.
 - At most one of **Patch**, **Centroid**, or **Sample** Decorations can be applied to the same object or member.
 - At most one of **RowMajor** and **ColMajor** Decorations can be applied to a structure type.
 - At most one of **Block** and **BufferBlock** Decorations can be applied to a structure type.
- All <id> used for Scope and Memory Semantics must be of an OpConstant.

2.16.3 Validation Rules for Kernel Capabilities

• The Signedness in **OpTypeInt** must always be 0.

2.17 Universal Limits

These quantities are minimum limits for all implementations and validators. Implementations are allowed to support larger quantities. Specific APIs may impose larger minimums. See Language Capabilities.

Validators must either

- inform when these limits are crossed, or
- be explicitly parameterized with larger limits.

Table 3: Limits

Limited Freiter	Minimum Limit		
Limited Entity	Decimal	Hexadecimal	
Characters in a literal string	65,535	FFFF	
Instruction word count	65,535	FFFF	
Result <id> bound</id>			
	4,194,303	3FFFFF	
See Physical Layout for the shader-specific bound.			
Control-flow nesting depth			
Measured per function, in program order, counting the maximum number of OpBranch, OpBranchConditional, or OpSwitch that are seen without yet seeing their corresponding <i>Merge Block</i> , as declared by OpSelectionMerge or OpLoopMerge.	1023	3FF	
Global variables (Storage Class other than Function)	65,535	FFFF	
Local variables (Function Storage Class)	524,287	7FFFF	
Decorations per target <id></id>		of entries in the ecoration table.	
Execution modes per entry point	255	FF	
Indexes for OpAccessChain, OpInBoundsAccessChain, OpPtrAccessChain, OpInBoundsPtrAccessChain, OpCompositeExtract, and OpCompositeInsert	255	FF	
Number of function parameters, per function declaration	255	FF	
OpFunctionCall actual arguments	255	FF	
OpExtInst actual arguments	255	FF	
OpSwitch (literal, label) pairs	16,383	3FFF	
OpTypeStruct members	16,383	3FFF	
Structure nesting depth	255	FF	

2.18 Memory Model

A memory model is chosen using a single OpMemoryModel instruction near the beginning of the module. This selects both an addressing model and a memory model.

The **Logical** addressing model means pointers are abstract, having no physical size or numeric value. In this mode, pointers can only be created from existing objects, and they cannot be stored into an object.

The non-**Logical** addressing models allow physical pointers to be formed. OpVariable can be used to create objects that hold pointers. These are declared for a specific Storage Class. Pointers for one Storage Class cannot be used to access

objects in another Storage Class. However, they can be converted with conversion opcodes. Any particular addressing model must describe the bit width of pointers for each of the storage classes.

2.18.1 Memory Layout

When memory is shared between a SPIR-V module and an API, its contents are transparent, and must be agreed on. For example, the **Offset**, **MatrixStride**, and **ArrayStride** Decorations applied to members of a struct object can partially define how the memory is laid out. In addition, the following are always true, applied recursively as needed, of the offsets within the memory buffer:

- a vector consumes contiguous memory with lower-numbered components appearing in smaller offsets than higher-numbered components, and with component 0 starting at the vector's **Offset** Decoration, if present
- in an array, lower-numbered elements appear at smaller offsets than higher-numbered elements, with element 0 starting at the **Offset** Decoration for the array, if present
- a structure has lower-numbered members appearing at smaller offsets than higher-numbered members, with member 0 starting at the **Offset** Decoration for the structure, if present
- in a matrix, lower-numbered columns appear at smaller offsets than higher-numbered columns, and lower-numbered components within the matrix's vectors appearing at smaller offsets than high-numbered components, with component 0 of column 0 starting at the **Offset** Decoration, if present (the **RowMajor** and **ColMajor** Decorations dictate what is contiguous)

2.18.2 Aliasing

Here, aliasing means one of:

- Two or more pointers that point into overlapping parts of the same underlying object. That is, two intermediates, both of which are typed pointers, that can be dereferenced (in bounds) such that both dereferences access the same memory.
- Images, buffers, or other externally allocated objects where a function might access the same underlying memory via accesses to two different objects.

How aliasing is managed depends on the Memory Model:

- The simple and GLSL memory models can assume that aliasing is generally not present. Specifically, the compiler is free to compile as if aliasing is not present, unless a pointer is explicitly indicated to be an alias. This is indicated by applying the **Aliased Decoration** to an *intermediate* object's <id>. Applying **Restrict** is allowed, but has no effect.
- The OpenCL memory models must assume that aliasing is generally present. Specifically, the compiler must compile as if aliasing is present, unless a pointer is explicitly indicated to not alias. This is done by applying the **Restrict** Decoration to an *intermediate* object's <id>. Applying **Aliased** is allowed, but has no effect.

It is invalid to apply both **Restrict** and **Aliased** to the same <*id*>.

2.19 Derivatives

Derivatives appear only in the **Fragment** Execution Model. They can be implicit or explicit. Some image instructions consume implicit derivatives, while the derivative instructions compute explicit derivatives. In all cases, derivatives are well defined only if the derivative group has uniform control flow.

2.20 Code Motion

Texturing instructions in the Fragment Execution Model that rely on an implicit derivative cannot be moved into control flow that is not known to be uniform control flow within each derivative group.

3 Binary Form

This section contains the exact form for all instructions, starting with the numerical values for all fields. See Physical Layout for the order words appear in.

3.1 Magic Number

Magic number for a SPIR-V module.

Tip

Endianness: A module is defined as a stream of words, not a stream of bytes. However, if stored as a stream of bytes (e.g., in a file), the magic number can be used to deduce what endianness to apply to convert the byte stream back to a word stream.

Magic Number	
0x07230203	

3.2 Source Language

The source language is for debug purposes only, with no semantics that affect the meaning of other parts of the module. Used by OpSource.

Source Language		
0	0 Unknown	
1	ESSL	
2	GLSL	
3	OpenCL_C	
4	OpenCL_CPP	

3.3 Execution Model

Used by OpEntryPoint.

	Execution Model	Required Capability
0	Vertex	Shader
	Vertex shading stage.	
1	TessellationControl	Tessellation
	Tessellation control (or hull) shading stage.	
2	TessellationEvaluation	Tessellation
	Tessellation evaluation (or domain) shading	
	stage.	
3	Geometry	Geometry
	Geometry shading stage.	
4	Fragment	Shader
	Fragment shading stage.	
5	GLCompute	Shader
	Graphical compute shading stage.	
6	Kernel	Kernel
	Compute kernel.	

3.4 Addressing Model

Used by OpMemoryModel.

Addressing Model		Required Capability
0	Logical	
1	Physical32	Addresses
	Indicates a 32-bit module, where the address	
	width is equal to 32 bits.	
2	Physical64	Addresses
	Indicates a 64-bit module, where the address	
	width is equal to 64 bits.	

3.5 Memory Model

Used by OpMemoryModel.

Memory Model		Required Capability
0	Simple	Shader
	No shared memory consistency issues.	
1	GLSL450	Shader
	Memory model needed by later versions of	
	GLSL and ESSL. Works across multiple	
	versions.	
2	OpenCL	Kernel
	OpenCL memory model.	

3.6 Execution Mode

Declare the modes an entry point will execute in. Used by OpExecutionMode.

Execution Mode		Required Capability	Extra Operands
0	Invocations	Geometry	Literal Number
	Number of times to invoke the geometry		Number of invocations
	stage for each input primitive received. The		
	default is to run once for each input		
	primitive. If greater than the		
	target-dependent maximum, it will fail to		
	compile. Only valid with the Geometry		
	Execution Model.		
1	SpacingEqual	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into a collection of		
	equal-sized segments. Only valid with one		
	of the tessellation Execution Models.		
2	SpacingFractionalEven	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into an even number of		
	equal-length segments plus two additional		
	shorter fractional segments. Only valid with		
	one of the tessellation Execution Models.		

	Execution Mode	Required Capability	Extra Operands
3	SpacingFractionalOdd	Tessellation	
	Requests the tessellation primitive generator		
	to divide edges into an odd number of		
	equal-length segments plus two additional		
	shorter fractional segments. Only valid with		
	one of the tessellation Execution Models.		
4	VertexOrderCw	Tessellation	
	Requests the tessellation primitive generator		
	to generate triangles in clockwise order.		
	Only valid with one of the tessellation		
	Execution Models.		
5	VertexOrderCcw	Tessellation	
	Requests the tessellation primitive generator		
	to generate triangles in counter-clockwise		
	order. Only valid with one of the		
6	tessellation Execution Models.	Shader	
0	PixelCenterInteger Pixels appear centered on whole-number	Snader	
	pixel offsets. E.g., the coordinate (0.5, 0.5)		
	appears to move to (0.0, 0.0). Only valid		
	with the Fragment Execution Model. If a		
	Fragment entry point does not have this		
	set, pixels appear centered at offsets of (0.5,		
	0.5) from whole numbers		
7	OriginUpperLeft	Shader	
	Pixel coordinates appear to originate in the		
	upper left, and increase toward the right and		
	downward. Only valid with the Fragment		
	Execution Model.		
8	OriginLowerLeft	Shader	
	Pixel coordinates appear to originate in the		
	lower left, and increase toward the right and		
	upward. Only valid with the Fragment		
	Execution Model.		
9	EarlyFragmentTests	Shader	
	Fragment tests are to be performed before		
	fragment shader execution. Only valid with		
10	the Fragment Execution Model.	Tonnell - 42	
10	PointMode Requests the taggellation mimitive generator	Tessellation	
	Requests the tessellation primitive generator		
	to generate a point for each distinct vertex		
	in the subdivided primitive, rather than to generate lines or triangles. Only valid with		
	one of the tessellation Execution Models.		
11	Xfb	TransformFeedback	
11	This stage will run in transform	Transformir coupack	
	feedback-capturing mode and this module is		
	responsible for describing the		
	transform-feedback setup. See the		
	XfbBuffer, Offset, and XfbStride		
	Decorations.		

	Execution Mode	Required Capability	Extra Op	erands	
12	DepthReplacing	Shader	•		
	This mode must be declared if this module				
	potentially changes the fragment's depth.				
	Only valid with the Fragment Execution				
	Model.				
14	DepthGreater	Shader			
	External optimizations may assume depth				
	modifications will leave the fragment's				
	depth as greater than or equal to the				
	fragment's interpolated depth value (given				
	by the z component of the FragCoord				
	BuiltIn decorated variable). Only valid with				
	the Fragment Execution Model.				
15	DepthLess	Shader			
	External optimizations may assume depth				
	modifications leave the fragment's depth				
	less than the fragment's interpolated depth				
	value, (given by the z component of the				
	FragCoord BuiltIn decorated variable).				
	Only valid with the Fragment Execution				
	Model.				
16	DepthUnchanged	Shader			
	External optimizations may assume this				
	stage did not modify the fragment's depth.				
	However, DepthReplacing mode must				
	accurately represent depth modification.				
	Only valid with the Fragment Execution				
	Model.				
17	LocalSize		Literal	Literal	Literal
	Indicates the work-group size in the x , y ,		Number	Number	Number
	and z dimensions. Only valid with the		x size	y size	z size
	GLCompute or Kernel Execution Models.				
18	LocalSizeHint	Kernel	Literal	Literal	Literal
	A hint to the compiler, which indicates the		Number	Number	Number
	most likely to be used work-group size in		x size	y size	z size
	the x , y , and z dimensions. Only valid with				
	the Kernel Execution Model.				
19	InputPoints	Geometry			
	Stage input primitive is <i>points</i> . Only valid				
	with the Geometry Execution Model.				
20	InputLines	Geometry			
	Stage input primitive is <i>lines</i> . Only valid				
21	with the Geometry Execution Model.				
21	InputLinesAdjacency	Geometry			
	Stage input primitive is <i>lines adjacency</i> .				
	Only valid with the Geometry Execution				
22	Model.	Comment The 11 49			
22	Triangles	Geometry, Tessellation			
	For a geometry stage, input primitive is				
	triangles. For a tessellation stage, requests				
	the tessellation primitive generator to				
	generate triangles. Only valid with the				
	Geometry or one of the tessellation				
	Execution Models.				

	Execution Mode	Required Capability	Extra Operands
23	InputTrianglesAdjacency	Geometry	_
	Geometry stage input primitive is <i>triangles</i>		
	adjacency. Only valid with the Geometry		
	Execution Model.		
24	Quads	Tessellation	
	Requests the tessellation primitive generator		
	to generate <i>quads</i> . Only valid with one of		
	the tessellation Execution Models.		
25	Isolines	Tessellation	
	Requests the tessellation primitive generator		
	to generate <i>isolines</i> . Only valid with one of		
26	the tessellation Execution Models.		Y' 1NY 1
26	OutputVertices	Geometry, Tessellation	Literal Number
	For a geometry stage, the maximum number		Vertex count
	of vertices the shader will ever emit in a		
	single invocation. For a tessellation-control		
	stage, the number of vertices in the output patch produced by the tessellation control		
	shader, which also specifies the number of		
	times the tessellation control shader is		
	invoked. Only valid with the Geometry or		
	one of the tessellation Execution Models.		
27	OutputPoints	Geometry	
2,	Stage output primitive is <i>points</i> . Only valid	Geometry	
	with the Geometry Execution Model.		
28	OutputLineStrip	Geometry	
	Stage output primitive is <i>line strip</i> . Only		
	valid with the Geometry Execution Model.		
29	OutputTriangleStrip	Geometry	
	Stage output primitive is <i>triangle strip</i> .	-	
	Only valid with the Geometry Execution		
	Model.		
30	VecTypeHint	Kernel	Literal Number
	A hint to the compiler, which indicates that		Vector type
	most operations used in the entry point are		
	explicitly vectorized using a particular		
	vector type. The 16 high-order bits of		
	Vector Type operand specify the number of		
	components of the vector. The 16 low-order		
	bits of <i>Vector Type</i> operand specify the <i>data</i>		
	type of the vector.		
	Those are the legal data two values.		
	These are the legal <i>data type</i> values: 0 represents an 8-bit integer value.		
	<i>I</i> represents a 16-bit integer value.		
	2 represents a 32-bit integer value.		
	3 represents a 64-bit integer value.		
	4 represents a 16-bit float value.		
	5 represents a 32-bit float value.		
	6 represents a 64-bit float value.		
	£		
	Only valid with the Kernel Execution		
	Model.		
	I	Ĭ.	I.

	Execution Mode	Required Capability	Extra Operands
31	ContractionOff	Kernel	
	Indicates that floating-point-expressions contraction is disallowed. Only valid with the Kernel Execution Model.		

3.7 Storage Class

Class of storage for declared variables (does not include intermediate values). Used by:

- OpTypePointer
- OpTypeForwardPointer
- OpVariable
- OpGenericCastToPtrExplicit

	Storage Class	Required Capability
0	UniformConstant	
	Shared externally, visible across all functions	
	in all invocations in all work groups.	
	Graphics uniform memory. OpenCL constant	
	memory. Read only.	
1	Input	Shader
	Input from pipeline. Visible across all	
	functions in the current invocation. Read	
	only.	
2	Uniform	Shader
	Shared externally, visible across all functions	
	in all invocations in all work groups.	
	Graphics uniform blocks and buffer blocks.	
3	Output	Shader
	Output to pipeline. Visible across all	
	functions in the current invocation.	
4	Workgroup	
	Shared across all invocations within a work	
	group. Visible across all functions. The	
	OpenGL "shared" storage qualifier. OpenCL	
	local memory.	
5	CrossWorkgroup	
	Visible across all functions of all invocations	
	of all work groups. OpenCL global memory.	
6	Private	Shader
	Visible to all functions in the current	
	invocation. Regular global memory.	
7	Function	
	Visible only within the declaring function of	
	the current invocation. Regular function	
	memory.	
8	Generic	GenericPointer
-	For generic pointers, which overload the	
	Function, Workgroup, and	
	CrossWorkgroup Storage Classes.	

	Storage Class	Required Capability
9	PushConstant	Shader
	For holding push-constant memory, visible	
	across all functions in all invocations in all	
	work groups. Read only. Intended to contain	
	a small bank of values pushed from the API.	
10	AtomicCounter	AtomicStorage
	For holding atomic counters. Visible across	
	all functions of the current invocation.	
	Atomic counter-specific memory.	
11	Image	
	For holding image memory.	

3.8 Dim

Dimensionality of an image. Used by OpTypeImage.

	Dim	Required Capability
0	1D	Sampled1D
1	2D	
2	3D	
3	Cube	Shader
4	Rect	SampledRect
5	Buffer	SampledBuffer
6	SubpassData	InputAttachment

3.9 Sampler Addressing Mode

Addressing mode for creating constant samplers. Used by OpConstantSampler.

	Sampler Addressing Mode	Required Capability
0	None	Kernel
	The image coordinates used to sample	
	elements of the image refer to a location	
	inside the image, otherwise the results are	
	undefined.	
1	ClampToEdge	Kernel
	Out-of-range image coordinates are clamped	
	to the extent.	
2	Clamp	Kernel
	Out-of-range image coordinates will return a	
	border color.	
3	Repeat	Kernel
	Out-of-range image coordinates are wrapped	
	to the valid range. Can only be used with	
	normalized coordinates.	
4	RepeatMirrored	Kernel
	Flip the image coordinate at every integer	
	junction. Can only be used with normalized	
	coordinates.	

3.10 Sampler Filter Mode

Filter mode for creating constant samplers. Used by OpConstantSampler.

	Sampler Filter Mode	Required Capability
0	Nearest	Kernel
	Use filter nearest mode when performing a	
	read image operation.	
1	Linear	Kernel
	Use filter linear mode when performing a	
	read image operation.	

3.11 Image Format

Declarative image format. Used by OpTypeImage.

	Image Format	Required Capability
0	Unknown	Trequired Supulmey
1	Rgba32f	Shader
2	Rgba16f	Shader
3	R32f	Shader
4	Rgba8	Shader
5	Rgba8Snorm	Shader
6	Rg32f	StorageImageExtendedFormats
7	Rg16f	StorageImageExtendedFormats
8	R11fG11fB10f	StorageImageExtendedFormats
9	R16f	StorageImageExtendedFormats
10	Rgba16	StorageImageExtendedFormats
11	Rgb10A2	StorageImageExtendedFormats
12	Rg16	StorageImageExtendedFormats
13	Rg8	StorageImageExtendedFormats
14	R16	StorageImageExtendedFormats
15	R8	StorageImageExtendedFormats
16	Rgba16Snorm	StorageImageExtendedFormats
17	Rg16Snorm	StorageImageExtendedFormats
18	Rg8Snorm	StorageImageExtendedFormats
19	R16Snorm	StorageImageExtendedFormats
20	R8Snorm	StorageImageExtendedFormats
21	Rgba32i	Shader
22	Rgba16i	Shader
23	Rgba8i	Shader
24	R32i	Shader
25	Rg32i	StorageImageExtendedFormats
26	Rg16i	StorageImageExtendedFormats
27	Rg8i	StorageImageExtendedFormats
28	R16i	StorageImageExtendedFormats
29	R8i	StorageImageExtendedFormats
30	Rgba32ui	Shader
31	Rgba16ui	Shader
32	Rgba8ui	Shader
33	R32ui	Shader
34	Rgb10a2ui	StorageImageExtendedFormats
35	Rg32ui	StorageImageExtendedFormats

	Image Format	Required Capability
36	Rg16ui	StorageImageExtendedFormats
37	Rg8ui	StorageImageExtendedFormats
38	R16ui	StorageImageExtendedFormats
39	R8ui	StorageImageExtendedFormats

3.12 Image Channel Order

Image channel order returned by OpImageQueryOrder.

	Image Channel Order	Required Capability
0	R	Kernel
1	A	Kernel
2	RG	Kernel
3	RA	Kernel
4	RGB	Kernel
5	RGBA	Kernel
6	BGRA	Kernel
7	ARGB	Kernel
8	Intensity	Kernel
9	Luminance	Kernel
10	Rx	Kernel
11	RGx	Kernel
12	RGBx	Kernel
13	Depth	Kernel
14	DepthStencil	Kernel
15	sRGB	Kernel
16	sRGBx	Kernel
17	sRGBA	Kernel
18	sBGRA	Kernel

3.13 Image Channel Data Type

Image channel data type returned by OpImageQueryFormat.

	Image Channel Data Type	Required Capability
0	SnormInt8	Kernel
1	SnormInt16	Kernel
2	UnormInt8	Kernel
3	UnormInt16	Kernel
4	UnormShort565	Kernel
5	UnormShort555	Kernel
6	UnormInt101010	Kernel
7	SignedInt8	Kernel
8	SignedInt16	Kernel
9	SignedInt32	Kernel
10	UnsignedInt8	Kernel
11	UnsignedInt16	Kernel
12	UnsignedInt32	Kernel
13	HalfFloat	Kernel
14	Float	Kernel
15	UnormInt24	Kernel

Image Channel Data Type		Required Capability
16	UnormInt101010_2	Kernel

3.14 Image Operands

Additional operands to sampling, or getting texels from, an image. Bits that are set can indicate that another operand follows. If there are multiple following operands indicated, they are ordered: Those indicated by smaller-numbered bits appear first. At least one bit must be set (**None** is invalid).

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

Used by:

- OpImageSampleImplicitLod
- OpImageSampleExplicitLod
- OpImageSampleDrefImplicitLod
- OpImageSampleDrefExplicitLod
- OpImageSampleProjImplicitLod
- OpImageSampleProjExplicitLod
- OpImageSampleProjDrefImplicitLod
- OpImageSampleProjDrefExplicitLod
- OpImageFetch
- OpImageGather
- OpImageDrefGather
- OpImageRead
- OpImageWrite
- OpImageSparseSampleImplicitLod
- OpImageSparseSampleExplicitLod
- OpImageSparseSampleDrefImplicitLod
- OpImageSparseSampleDrefExplicitLod
- OpImageSparseSampleProjImplicitLod
- OpImageSparseSampleProjExplicitLod
- OpImageSparseSampleProjDrefImplicitLod
- OpImageSparseSampleProjDrefExplicitLod
- OpImageSparseFetch
- OpImageSparseGather
- OpImageSparseDrefGather
- OpImageSparseRead

	Image Operands	Required Capability
0x0	None	
0x1	Bias	Shader
	A following operand is the bias added to	
	the implicit level of detail. Only valid with	
	implicit-lod instructions. It must be a	
	floating-point type scalar. This can only be	
	used with an OpTypeImage that has a Dim	
	operand of 1D, 2D, 3D, or Cube, and the	
	MS operand must be 0.	

	Image Operands	Required Capability
0x2	Lod	
	A following operand is the explicit	
	level-of-detail to use. Only valid with	
	explicit-lod instructions. For sampling	
	operations, it must be a floating-point type	
	scalar. For queries and fetch operations, it	
	must be an integer type scalar. This can	
	only be used with an OpTypeImage that	
	has a Dim operand of 1D, 2D, 3D, or	
	Cube , and the <i>MS</i> operand must be 0.	
0x4	Grad	
OA I	Two following operands are dx followed	
	by dy . These are explicit derivatives in the	
	x and y direction to use in computing level	
	of detail. Each is a scalar or vector	
	containing $(du/dx[, dv/dx] [, dw/dx])$ and	
	(du/dy[, dv/dy] [, dw/dy]). The number of	
	components of each must equal the	
	number of components in <i>Coordinate</i> ,	
	minus the <i>array layer</i> component, if	
	present. Only valid with explicit-lod	
	instructions. They must be a scalar or	
	· · · · · · · · · · · · · · · · · · ·	
	vector of floating-point type. This can only	
	be used with an OpTypeImage that has an	
	MS operand of 0. It is invalid to set both	
0.0	the Lod and Grad bits.	
0x8	ConstOffset	
	A following operand is added to (u, v, w)	
	before texel lookup. It must be an $\langle id \rangle$ of	
	an integer-based constant instruction of	
	scalar or vector type. It is a compile-time	
	error if these fall outside a	
	target-dependent allowed range. The	
	number of components must equal the	
	number of components in <i>Coordinate</i> ,	
	minus the array layer component, if	
	present.	
0x10	Offset	ImageGatherExtended
	A following operand is added to (u, v, w)	
	before texel lookup. It must be a scalar or	
	vector of integer type. It is a compile-time	
	error if these fall outside a	
	target-dependent allowed range. The	
	number of components must equal the	
	number of components in Coordinate,	
	minus the array layer component, if	
	present.	

	Image Operands	Required Capability
0x20	ConstOffsets	
	A following operand is Offsets. Offsets	
	must be an <id> of a constant instruction</id>	
	making an array of size four of vectors of	
	two integer components. Each gathered	
	texel is identified by adding one of these	
	array elements to the (u, v) sampled	
	location. It is a compile-time error if this	
	falls outside a target-dependent allowed	
	range. Only valid with OpImageGather or	
	OpImageDrefGather.	
0x40	Sample	
	A following operand is the sample number	
	of the sample to use. Only valid with	
	OpImageFetch, OpImageRead, and	
	OpImageWrite. It is invalid to have a	
	Sample operand if the underlying	
	OpTypeImage has MS of 0. It must be an	
	integer type scalar.	
0x80	MinLod	MinLod
	A following operand is the minimum	
	level-of-detail to use when accessing the	
	image. Only valid with Implicit	
	instructions and Grad instructions. It must	
	be a floating-point type scalar. This can	
	only be used with an OpTypeImage that	
	has a Dim operand of 1D, 2D, 3D, or	
	Cube , and the <i>MS</i> operand must be 0.	

3.15 FP Fast Math Mode

Enables fast math operations which are otherwise unsafe.

• Only valid on OpFAdd, OpFSub, OpFMul, OpFDiv, OpFRem, and OpFMod instructions.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

FP Fast Math Mode		Required Capability
0x0	None	
0x1	NotNaN	Kernel
	Assume parameters and result are not	
	NaN.	
0x2	NotInf	Kernel
	Assume parameters and result are not +/-	
	Inf.	
0x4	NSZ	Kernel
	Treat the sign of a zero parameter or result	
	as insignificant.	
0x8	AllowRecip	Kernel
	Allow the usage of reciprocal rather than	
	perform a division.	

	FP Fast Math Mode	Required Capability
0x10	Fast	Kernel
	Allow algebraic transformations according	
	to real-number associative and distributive	
	algebra. This flag implies all the others.	

3.16 FP Rounding Mode

Associate a rounding mode to a floating-point conversion instruction.

By default

- Conversions from floating-point to integer types use the round-toward-zero rounding mode.
- Conversions to floating-point types use the round-to-nearest-even rounding mode.

	FP Rounding Mode	Required Capability
0	RTE	Kernel
	Round to nearest even.	
1	RTZ	Kernel
	Round towards zero.	
2	RTP	Kernel
	Round towards positive infinity.	
3	RTN	Kernel
	Round towards negative infinity.	

3.17 Linkage Type

Associate a linkage type to functions or global variables. See linkage.

Linkage Type		Required Capability	
0	Export	Linkage	
	Accessible by other modules as well.		
1	Import	Linkage	
	A declaration of a global variable or a		
	function that exists in another module.		

3.18 Access Qualifier

Defines the access permissions.

Used by OpTypeImage and OpTypePipe.

Access Qualifier		Required Capability
0	ReadOnly	Kernel
	A read-only object.	
1	WriteOnly	Kernel
	A write-only object.	
2	ReadWrite	Kernel
	A readable and writable object.	

3.19 Function Parameter Attribute

Adds additional information to the return type and to each parameter of a function.

	Function Parameter Attribute	Required Capability
0	Zext	Kernel
	Value should be zero extended if needed.	
1	Sext	Kernel
	Value should be sign extended if needed.	
2	ByVal	Kernel
	This indicates that the pointer parameter	
	should really be passed by value to the	
	function. Only valid for pointer parameters	
	(not for ret value).	
3	Sret	Kernel
	Indicates that the pointer parameter specifies	
	the address of a structure that is the return	
	value of the function in the source program.	
	Only applicable to the first parameter which	
	must be a pointer parameters.	
4	NoAlias	Kernel
	Indicates that the memory pointed by a	
	pointer parameter is not accessed via pointer	
	values which are not derived from this	
	pointer parameter. Only valid for pointer	
	parameters. Not valid on return values.	T7 1
5	NoCapture	Kernel
	The callee does not make a copy of the	
	pointer parameter into a location that is	
	accessible after returning from the callee.	
	Only valid for pointer parameters. Not valid on return values.	
6	on return values. NoWrite	Kernel
U		Kernei
	Can only read the memory pointed by a	
	pointer parameter. Only valid for pointer parameters. Not valid on return values.	
7	NoReadWrite	Kernel
/	Cannot dereference the memory pointed by a	Kerner
	pointer parameter. Only valid for pointer	
	parameters. Not valid on return values.	
	parameters. Not vand on return vardes.	

3.20 Decoration

Used by OpDecorate and OpMemberDecorate.

	Decoration	Required Capability	Extra Operands
0	RelaxedPrecision	Shader	
	Allow reduced precision operations. To be used		
	as described in Relaxed Precision.		
1	SpecId	Shader	Literal Number
	Apply to a scalar specialization constant. Forms		Specialization
	the API linkage for setting a specialized value.		Constant ID
	See specialization.		

	Decoration	Required Capability	Extra Operands
2	Block	Shader	
	Apply to a structure type to establish it is a		
	non-SSBO-like shader-interface block.		
3	BufferBlock	Shader	
	Apply to a structure type to establish it is an		
	SSBO-like shader-interface block.		
4	RowMajor	Matrix	
	Applies only to a member of a structure type.		
	Only valid on a matrix or array whose most basic		
	element is a matrix. Indicates that components		
	within a row are contiguous in memory.		
5	ColMajor	Matrix	
3	Applies only to a member of a structure type.	IVIUUI IX	
	Only valid on a matrix or array whose most basic		
	element is a matrix. Indicates that components		
	within a column are contiguous in memory.		
6	•	Shader	Literal Number
O	ArrayStride	Shauer	
	Apply to an array type to specify the stride, in		Array Stride
	bytes, of the array's elements. Must not be		
	applied to anything other than an array type.	37	
7	MatrixStride	Matrix	Literal Number
	Applies only to a member of a structure type.		Matrix Stride
	Only valid on a matrix or array whose most basic		
	element is a matrix. Specifies the stride of rows in		
	a RowMajor -decorated matrix, or columns in a		
	ColMajor-decorated matrix.		
8	GLSLShared	Shader	
	Apply to a structure type to get GLSL shared		
	memory layout.		
9	GLSLPacked	Shader	
	Apply to a structure type to get GLSL packed		
	memory layout.		
10	CPacked	Kernel	
	Apply to a structure type, to marks it as "packed",		
	indicating that the alignment of the structure is		
	one and that there is no padding between structure		
	members.		
11	BuiltIn		Literal Number
	Apply to an object or a member of a structure		See BuiltIn
	type. Indicates which built-in variable the entity		
	represents. See BuiltIn for more information.		
13	NoPerspective	Shader	
-	Apply to an object or a member of a structure		
	type. Indicates that linear, non-perspective		
	correct, interpolation must be used. The object or		
	member must be a scalar or vector of		
	floating-point type. Arrays of these types are also		
	allowed. Only valid for the Input and Output		
	Storage Classes.		
	Diorage Classes.		

	Decoration	Required Capability	Extra Operands
14	Flat	Shader	_
	Apply to an object or a member of a structure		
	type. Indicates no interpolation will be done. The		
	non-interpolated value will come from a vertex,		
	as described in the API specification. The object		
	or member must be a scalar or vector of		
	floating-point type or integer type. Arrays of		
	these types are also allowed. Only valid for the		
	Input and Output Storage Classes.		
15	Patch	Tessellation	
	Apply to an object or a member of a structure		
	type. Indicates a tessellation patch. The object or		
	member must be a scalar or vector of		
	floating-point type. Arrays of these types are also		
	allowed. Only valid for the Input and Output		
	Storage Classes. Invalid to use on objects or types		
1.6	referenced by non-tessellation Execution Models.	GI I	
16	Centroid	Shader	
	Apply to an object or a member of a structure		
	type. When used with multi-sampling		
	rasterization, allows a single interpolation		
	location for an entire pixel. The interpolation location must lie in both the pixel and in the		
	primitive being rasterized. The object or member		
	must be a scalar or vector of floating-point type.		
	Arrays of these types are also allowed. Only valid		
	for the Input and Output Storage Classes.		
17	Sample	SampleRateShading	
1,	Apply to an object or a member of a structure	Sumpressuresnaumg	
	type. When used with multi-sampling		
	rasterization, requires per-sample interpolation.		
	The interpolation locations must be the locations		
	of the samples lying in both the pixel and in the		
	primitive being rasterized. The object or member		
	must be a scalar or vector of floating-point type.		
	Arrays of these types are also allowed. Only valid		
	for the Input and Output Storage Classes.		
18	Invariant	Shader	
	Apply to a variable, to indicate expressions		
	computing its value be done invariant with respect		
	to other modules computing the same		
	expressions.		
19	Restrict		
	Apply to a variable, to indicate the compiler may		
	compile as if there is no aliasing. See the Aliasing		
20	section for more detail.		
20	Aliased		
	Apply to a variable, to indicate the compiler is to		
	generate accesses to the variable that work		
	correctly in the presence of aliasing. See the		
	Aliasing section for more detail.		

	Decoration	Required Capability	Extra Operands
21	Volatile		
	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the		
	Uniform Storage Class with the BufferBlock		
	Decoration. This indicates the memory holding		
	the variable is volatile memory. Accesses to		
	volatile memory cannot be eliminated, duplicated,		
	or combined with other accesses. The variable		
22	cannot be in the Function Storage Class.	Variation	
22	Constant Indicates that a glabal wariable is constant and	Kernel	
	Indicates that a global variable is constant and will never be modified. Only allowed on global		
	variables.		
23	Coherent		
2.5	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the		
	Uniform Storage Class with the BufferBlock		
	Decoration. This indicates the memory backing		
	the object is coherent.		
24	NonWritable		
	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the		
	Uniform Storage Class with the BufferBlock		
	Decoration. This indicates the memory holding		
	the variable is not writable, and that this module		
	does not write to it.		
25	NonReadable		
	Apply to an object or a member of a structure		
	type. Can only be used for objects declared as		
	storage images (see OpTypeImage) or in the Uniform Storage Class with the BufferBlock		
	Decoration. This indicates the memory holding		
	the variable is not readable, and that this module		
	does not read from it.		
26	Uniform	Shader	
	Apply to an object or a member of a structure	22-2-	
	type. Asserts that the value backing the decorated		
	<id> is dynamically uniform, hence the</id>		
	consumer is allowed to assume this is the case.		
28	SaturatedConversion	Kernel	
	Indicates that a conversion to an integer type		
	which is outside the representable range of <i>Result</i>		
	Type will be clamped to the nearest representable		
	value of <i>Result Type</i> . <i>NaN</i> will be converted to 0.		
	This decoration can only be applied to conversion		
	instructions to integer types, not including the		
	OpSatConvertUToS and OpSatConvertSToU		
	instructions.		

	Decoration	Required Capability	Extra Operands
29	Stream	GeometryStreams	Literal Number
	Apply to an object or a member of a structure		Stream Number
	type. Indicates the stream number to put an		
	output on. Only valid for the Output Storage		
20	Class and the Geometry Execution Model.	GL 1	T. 137 1
30	Location	Shader	Literal Number
	Apply to a variable or a structure-type member.		Location
	Forms the main linkage for Storage Class Input		
	and Output variables:		
	- between the API and vertex-stage inputs,		
	- between consecutive programmable stages, or		
	- between fragment-stage outputs and the API.		
	Also can tag variables or structure-type members		
	in the UniformConstant Storage Class for		
	linkage with the API.		
	Only valid for the Input , Output , and		
31	UniformConstant Storage Classes.	Shader	Literal Number
31	Apply to an object or a member of a structure	Shader	
	type. Indicates which component within a		Component
	Location will be taken by the decorated entity.		
	Only valid for the Input and Output Storage		
	Classes.		
32	Index	Shader	Literal Number
-	Apply to a variable to identify a blend equation	21111112	Index
	input index, used as described in the API		
	specification. Only valid for the Output Storage		
	Class and the Fragment Execution Model.		
33	Binding	Shader	Literal Number
	Apply to a variable. Part of the main linkage		Binding Point
	between the API and SPIR-V modules for		
	memory buffers, images, etc. See the API		
	specification for more information.		
34	DescriptorSet	Shader	Literal Number
	Apply to a variable. Part of the main linkage		Descriptor Set
	between the API and SPIR-V modules for		
	memory buffers, images, etc. See the API		
	specification for more information.		
35	Offset	Shader	Literal Number
	Apply to a structure-type member. This gives the		Byte Offset
	byte offset of the member relative to the		
	beginning of the structure. Can be used, for		
	example, by both uniform and		
	transform-feedback buffers. It must not cause any		
	overlap of the structure's members, or overflow of		
	a transform-feedback buffer's XfbStride .		
36	XfbBuffer	TransformFeedback	Literal Number
	Apply to an object or a member of a structure		XFB Buffer Number
	type. Indicates which transform-feedback buffer		
	an output is written to. Only valid for the Output		
	Storage Classes of vertex processing Execution		
	Models.		

	Decoration	Required Capability	Extra Operands
37	XfbStride	TransformFeedback	Literal Number
	Apply to anything XfbBuffer is applied to.		XFB Stride
	Specifies the stride, in bytes, of		
	transform-feedback buffer vertices. If the		
	transform-feedback buffer is capturing any		
	double-precision components, the stride must be a		
	multiple of 8, otherwise it must be a multiple of 4.		
38	FuncParamAttr	Kernel	Function Parameter
	Indicates a function return value or parameter		Attribute
	attribute.		Function Parameter
			Attribute
39	FP Rounding Mode	Kernel	FP Rounding Mode
	Indicates a floating-point rounding mode.		Floating-Point
			Rounding Mode
40	FP Fast Math Mode	Kernel	FP Fast Math Mode
	Indicates a floating-point fast math flag.		Fast-Math Mode
41	Linkage Attributes	Linkage	Literal Linkage
	Associate linkage attributes to values. Only valid	_	String Type
	on OpFunction or global (module scope)		Name Linkage
	OpVariable. See linkage.		Type
42	NoContraction	Shader	1
	Apply to an arithmetic instruction to indicate the		
	operation cannot be combined with another		
	instruction to form a single operation. For		
	example, if applied to an OpFMul, that multiply		
	can't be combined with an addition to yield a		
	fused multiply-add operation. Furthermore, such		
	operations are not allowed to reassociate; e.g.,		
	add(a + add(b+c)) cannot be transformed to		
	add(add(a+b) + c).		
43	InputAttachmentIndex	InputAttachment	Literal Number
	Apply to a variable to provide an input-target		Attachment Index
	index (as described in the API specification).		
	Only valid in the Fragment Execution Model and		
	for variables of type OpTypeImage with a Dim		
	operand of SubpassData.		
44	Alignment	Kernel	Literal Number
	Apply to a pointer. This declares a known		Alignment
	minimum alignment the pointer has.		

3.21 BuiltIn

Used when Decoration is BuiltIn. Apply to either

- the result $\langle id \rangle$ of the variable declaration of the built-in variable, or
- a structure-type member, if the built-in is a member of a structure.

As stated per entry below, these have additional semantics and constraints described by the client API.

	BuiltIn	Required Capability
0	Position	Shader
	Output vertex position from a vertex	
	processing Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	
1	PointSize	Shader
	Output point size from a vertex processing	
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
3	ClipDistance	ClipDistance
	Array of clip distances. See Vulkan or	-
	OpenGL API specifications for more detail.	
4	CullDistance	CullDistance
	Array of clip distances. See Vulkan or	
	OpenGL API specifications for more detail.	
5	VertexId	Shader
	Input vertex ID to a Vertex Execution	
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
6	InstanceId	Shader
	Input instance ID to a Vertex Execution	2
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
7	PrimitiveId	Geometry, Tessellation
,	Primitive ID in a Geometry Execution	geometry, ressentation
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
8	InvocationId	Geometry, Tessellation
	Invocation ID, input to Geometry and	Geometry, ressentation
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
9	Layer	Geometry
	Layer output by a Geometry Execution	Jestines, j
	Model, input to a Fragment Execution	
	Model, for multi-layer framebuffer. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
10	ViewportIndex	MultiViewport
	Viewport Index output by a Geometry stage,	
	input to a Fragment Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
11	TessLevelOuter	Tessellation
	Output patch outer levels in a	
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
12	TessLevelInner	Tessellation
	Output patch inner levels in a	
	TessellationControl Execution Model. See	
	Vulkan or OpenGL API specifications for	
	more detail.	
	I and the second	I.

	BuiltIn	Required Capability
13	TessCoord	Tessellation
	Input vertex position in	
	TessellationEvaluation Execution Model.	
	See Vulkan or OpenGL API specifications	
	for more detail.	
14	PatchVertices	Tessellation
	Input patch vertex count in a tessellation	
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
15	FragCoord	Shader
	Coordinates $(x, y, z, 1/w)$ of the current	
	fragment, input to the Fragment Execution	
	Model. See Vulkan or OpenGL API	
	specifications for more detail.	
16	PointCoord	Shader
	Coordinates within a <i>point</i> , input to the	
	Fragment Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	
17	FrontFacing	Shader
	Face direction, input to the Fragment	
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
18	SampleId	SampleRateShading
	Input sample number to the Fragment	•
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
19	SamplePosition	SampleRateShading
	Input sample position to the Fragment	1
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
20	SampleMask	SampleRateShading
	Input or output sample mask to the	•
	Fragment Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	
22	FragDepth	Shader
	Output fragment depth from the Fragment	
	Execution Model. See Vulkan or OpenGL	
	API specifications for more detail.	
23	HelperInvocation	Shader
	Input whether a helper invocation, to the	
	Fragment Execution Model. See Vulkan or	
	OpenGL API specifications for more detail.	
24	NumWorkgroups	
	Number of workgroups in GLCompute or	
	Kernel Execution Models. See OpenCL,	
	Vulkan, or OpenGL API specifications for	
	more detail.	
25	WorkgroupSize	
	Work-group size in GLCompute or Kernel	
	Execution Models. See OpenCL, Vulkan, or	
	OpenGL API specifications for more detail.	

	BuiltIn	Required Capability
26	WorkgroupId	• • •
	Work-group ID in GLCompute or Kernel	
	Execution Models. See OpenCL, Vulkan, or	
	OpenGL API specifications for more detail.	
27	LocalInvocationId	
	Local invocation ID in GLCompute or	
	Kernel Execution Models. See OpenCL,	
	Vulkan, or OpenGL API specifications for	
	more detail.	
28	GlobalInvocationId	
	Global invocation ID in GLCompute or	
	Kernel Execution Models. See OpenCL,	
	Vulkan, or OpenGL API specifications for	
	more detail.	
29	LocalInvocationIndex	
	Local invocation index in GLCompute	
	Execution Models. See Vulkan or OpenGL	
	API specifications for more detail.	
	Transfer and the second	
	Work-group Linear ID in Kernel Execution	
	Models. See OpenCL API specification for	
	more detail.	
30	WorkDim	Kernel
	Work dimensions in Kernel Execution	
	Models. See OpenCL API specification for	
	more detail.	
31	GlobalSize	Kernel
31	Global size in Kernel Execution Models. See	Tion not
	OpenCL API specification for more detail.	
32	EnqueuedWorkgroupSize	Kernel
02	Enqueued work-group size in Kernel	2202.102
	Execution Models. See OpenCL API	
	specification for more detail.	
33	GlobalOffset	Kernel
	Global offset in Kernel Execution Models.	2202.1102
	See OpenCL API specification for more	
	detail.	
34	GlobalLinearId	Kernel
	Global linear ID in Kernel Execution	
	Models. See OpenCL API specification for	
	more detail.	
36	SubgroupSize	Kernel
	Subgroup size in Kernel Execution Models.	
	See OpenCL API specification for more	
	detail.	
37	SubgroupMaxSize	Kernel
	Subgroup maximum size in Kernel	
	Execution Models. See OpenCL API	
	specification for more detail.	
38	NumSubgroups	Kernel
	Number of subgroups in Kernel Execution	THE THE
	Models. See OpenCL API specification for	
	more detail.	
	more detain.	

	BuiltIn	Required Capability
39	NumEnqueuedSubgroups	Kernel
	Number of enqueued subgroups in Kernel	
	Execution Models. See OpenCL API	
	specification for more detail.	
40	SubgroupId	Kernel
	Subgroup ID in Kernel Execution Models.	
	See OpenCL API specification for more	
	detail.	
41	SubgroupLocalInvocationId	Kernel
	Subgroup local invocation ID in Kernel	
	Execution Models. See OpenCL API	
	specification for more detail.	
42	VertexIndex	Shader
	Vertex index. See Vulkan or OpenGL API	
	specifications for more detail.	
43	InstanceIndex	Shader
	Instance index. See Vulkan or OpenGL API	
	specifications for more detail.	

3.22 Selection Control

This value is a mask; it can be formed by combining the bits from multiple rows in the table below. Used by OpSelectionMerge.

Selection Control		
0x0	None	
0x1	Flatten	
	Strong request, to the extent possible, to	
	remove the control flow for this selection.	
0x2	DontFlatten	
	Strong request, to the extent possible, to	
	keep this selection as control flow.	

3.23 Loop Control

This value is a mask; it can be formed by combining the bits from multiple rows in the table below. Used by OpLoopMerge.

Loop Control		
0x0	None	
0x1	Unroll	
	Strong request, to the extent possible, to	
	unroll or unwind this loop.	
0x2	DontUnroll	
	Strong request, to the extent possible, to	
	keep this loop as a loop, without unrolling.	

3.24 Function Control

This value is a mask; it can be formed by combining the bits from multiple rows in the table below. Used by OpFunction.

Function Control		
0x0	None	
0x1	Inline	
	Strong request, to the extent possible, to	
	inline the function.	
0x2	DontInline	
	Strong request, to the extent possible, to not	
	inline the function.	
0x4	Pure	
	Compiler can assume this function has no	
	side effect, but might read global memory	
	or read through dereferenced function	
	parameters. Always computes the same	
	result for the same argument values.	
0x8	Const	
	Compiler can assume this function has no	
	side effects, and will not access global	
	memory or dereference function	
	parameters. Always computes the same	
	result for the same argument values.	

3.25 Memory Semantics <id>

Must be an <id> of a 32-bit integer scalar that contains a mask. The rest of this description is about that mask.

Memory semantics define memory-order constraints, and on what storage classes those constraints apply to. The memory order constrains the allowed orders in which memory operations in this invocation can made visible to another invocation. The storage classes specify to which subsets of memory these constraints are to be applied. Storage classes not selected are not being constrained.

Despite being a mask and allowing multiple bits to be combined, at most one of the first four (low-order) bits can be set. Requesting both **Acquire** and **Release** semantics is done by setting the **AcquireRelease** bit, not by setting two bits.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

Used by:

- OpControlBarrier
- OpMemoryBarrier
- OpAtomicLoad
- OpAtomicStore
- OpAtomicExchange
- OpAtomicCompareExchange
- OpAtomicCompareExchangeWeak
- OpAtomicIIncrement
- OpAtomicIDecrement
- · OpAtomicIAdd

- OpAtomicISub
- OpAtomicSMin
- OpAtomicUMin
- OpAtomicSMax
- OpAtomicUMax
- OpAtomicAnd
- OpAtomicOr
- OpAtomicXor
- OpAtomicFlagTestAndSet
- OpAtomicFlagClear

	Memory Semantics	Required Capability
0x0	None (Relaxed)	
0x2	Acquire	
	All memory operations provided in	
	program order after this memory operation	
	will execute after this memory operation.	
0x4	Release	
	All memory operations provided in	
	program order before this memory	
	operation will execute before this memory	
	operation.	
0x8	AcquireRelease	
	Has the properties of both Acquire and	
	Release semantics. It is used for	
	read-modify-write operations.	
0x10	SequentiallyConsistent	
	All observers will see this memory access	
	in the same order with respect to other	
	sequentially-consistent memory accesses	
	from this invocation.	
0x40	UniformMemory	Shader
	Apply the memory-ordering constraints to	
	Uniform Storage Class memory.	
0x80	SubgroupMemory	
	Apply the memory-ordering constraints to	
	subgroup memory.	
0x100	WorkgroupMemory	
	Apply the memory-ordering constraints to	
	Workgroup Storage Class memory.	
0x200	CrossWorkgroupMemory	
	Apply the memory-ordering constraints to	
	CrossWorkgroup Storage Class memory.	
0x400	AtomicCounterMemory	AtomicStorage
	Apply the memory-ordering constraints to	
	AtomicCounter Storage Class memory.	
0x800	ImageMemory	
	Apply the memory-ordering constraints to	
	image contents (types declared by	
	OpTypeImage), or to accesses done	
	through pointers to the Image Storage	
	Class.	

3.26 Memory Access

Memory access semantics.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

Used by:

- OpLoad
- OpStore
- OpCopyMemory
- OpCopyMemorySized

Memory Access	
0x0	None
0x1	Volatile
	This access cannot be eliminated,
	duplicated, or combined with other
	accesses.
0x2	Aligned
	This access has a known alignment,
	provided as a literal in the next operand.
0x4	Nontemporal
	Hints that the accessed address is not likely
	to be accessed again in the near future.

3.27 Scope <id>

Must be an <id> of a 32-bit integer scalar that contains a mask. The rest of this description is about that mask.

The execution scope or memory scope of an operation. When used as a memory scope, it specifies the distance of synchronization from the current invocation. When used as an execution scope, it specifies the set of executing invocations taking part in the operation. Used by:

- OpControlBarrier
- OpMemoryBarrier
- OpAtomicLoad
- OpAtomicStore
- OpAtomicExchange
- OpAtomicCompareExchange
- OpAtomicCompareExchangeWeak
- OpAtomicIIncrement
- OpAtomicIDecrement
- OpAtomicIAdd
- OpAtomicISub
- OpAtomicSMin
- OpAtomicUMin
- OpAtomicSMax
- OpAtomicUMax

- OpAtomicAnd
- OpAtomicOr
- OpAtomicXor
- OpGroupAsyncCopy
- OpGroupWaitEvents
- OpGroupAll
- OpGroupAny
- OpGroupBroadcast
- OpGroupIAdd
- OpGroupFAdd
- OpGroupFMin
- OpGroupUMin
- OpGroupSMin
- OpGroupFMax
- OpGroupUMax
- OpGroupSMax
- OpGroupReserveReadPipePackets
- OpGroupReserveWritePipePackets
- OpGroupCommitReadPipe
- OpGroupCommitWritePipe
- OpAtomicFlagTestAndSet
- OpAtomicFlagClear

Scope		
0	CrossDevice	
	Scope crosses multiple devices.	
1	Device	
	Scope is the current device.	
2	Workgroup	
	Scope is the current workgroup.	
3	Subgroup	
	Scope is the current subgroup.	
4	Invocation	
	Scope is the current Invocation.	

3.28 Group Operation

Defines the class of workgroup or subgroup operation. Used by:

- OpGroupIAdd
- OpGroupFAdd
- OpGroupFMin
- OpGroupUMin
- OpGroupSMin

- OpGroupFMax
- OpGroupUMax
- OpGroupSMax

	Group Operation	Required Capability
0	Reduce	Kernel
	A reduction operation for all values of a	
	specific value X specified by invocations	
	within a workgroup.	
1	InclusiveScan	Kernel
	A binary operation with an identity I and n	
	(where n is the size of the workgroup)	
	elements[$a_0, a_1, \ldots a_{n-1}$] resulting in [$a_0, (a_0)$	
	op a_1), $(a_0 \text{ op } a_1 \text{ op } \text{ op } a_{n-1})$]	
2	ExclusiveScan	Kernel
	A binary operation with an identity I and n	
	(where n is the size of the workgroup)	
	elements[$a_0, a_1, \ldots a_{n-1}$] resulting in [I, a_0, \ldots]	
	$(a_0 \text{ op } a_1), \ldots (a_0 \text{ op } a_1 \text{ op } \ldots \text{ op } a_{n-2})].$	

3.29 Kernel Enqueue Flags

Specify when the child kernel begins execution.

Note: Implementations are not required to honor this flag. Implementations may not schedule kernel launch earlier than the point specified by this flag, however. Used by OpEnqueueKernel.

	Kernel Enqueue Flags	Required Capability
0	NoWait Indicates that the enqueued kernels do not need to wait for the parent kernel to finish execution before they begin execution.	Kernel
1	WaitKernel Indicates that all work-items of the parent kernel must finish executing and all immediate side effects committed before the enqueued child kernel may begin execution. Note: Immediate meaning not side effects resulting from child kernels. The side effects would include stores to global memory and pipe reads and writes.	Kernel
2	WaitWorkGroup Indicates that the enqueued kernels wait only for the workgroup that enqueued the kernels to finish before they begin execution. Note: This acts as a memory synchronization point between work-items in a work-group and child kernels enqueued by work-items in the work-group.	Kernel

3.30 Kernel Profiling Info

Specify the profiling information to be queried. Used by OpCaptureEventProfilingInfo.

This value is a mask; it can be formed by combining the bits from multiple rows in the table below.

	Kernel Profiling Info	Required Capability
0x0	None	
0x1	CmdExecTime	Kernel
	Indicates that the profiling info queried is	
	the execution time.	

3.31 Capability

Capabilities a module can declare it uses. All used capabilities must be declared, either directly or through a dependency: all capabilities that a declared capability depends on are automatically implied.

The **Depends On** column lists the dependencies for each capability. These are the ones implicitly declared. It is not necessary (but allowed) to declare a dependency for a declared capability.

See the capabilities section for more detail. Used by OpCapability.

	Capability	Depends On
0	Matrix	
	Uses OpTypeMatrix.	
1	Shader	Matrix
	Uses Vertex, Fragment, or GLCompute	
	Execution Models.	
2	Geometry	Shader
	Uses the Geometry Execution Model.	
3	Tessellation	Shader
	Uses the TessellationControl or	
	TessellationEvaluation Execution Models.	
4	Addresses	
	Uses physical addressing, non-logical	
	addressing modes.	
5	Linkage	
	Uses partially linked modules and libraries.	
6	Kernel	
	Uses the Kernel Execution Model.	
7	Vector16	Kernel
	Uses OpTypeVector to declare 8 component	
	or 16 component vectors.	
8	Float16Buffer	Kernel
	Uses pointers to 16-bit floating-point data	
	types (but doesn't use 16-bit OpTypeFloat as	
	a Result <id>).</id>	
9	Float16	
	Uses OpTypeFloat to declare the 16-bit	
	floating-point type.	
10	Float64	
	Uses OpTypeFloat to declare the 64-bit	
	floating-point type.	

	Capability	Depends On
11	Int64	•
	Uses OpTypeInt to declare 64-bit integer	
	types.	
12	Int64Atomics	Int64
	Uses atomic instructions on 64-bit integer	
	types.	
13	ImageBasic	Kernel
	Uses OpTypeImage or OpTypeSampler in a	
	Kernel.	
14	ImageReadWrite	ImageBasic
	Uses OpTypeImage with the ReadWrite	
	access qualifier.	
15	ImageMipmap	ImageBasic
	Uses non-zero Lod Image Operands.	
17	Pipes	Kernel
	Uses OpTypePipe, OpTypeReserveId, or pipe	
	instructions.	
18	Groups	
	Uses group instructions.	
19	DeviceEnqueue	Kernel
	Uses OpTypeQueue, OpTypeDeviceEvent,	
20	and device side enqueue instructions.	
20	LiteralSampler	Kernel
	Samplers are made from literals within the	
21	module. See OpConstantSampler.	
21	AtomicStorage	Shader
22	Uses the AtomicCounter Storage Class.	
22	Int16 Uses OnTangint to dealers 16 hit integer	
	Uses OpTypeInt to declare 16-bit integer types.	
23	TessellationPointSize	Tessellation
23	Tessellation stage exports point size.	ressenation
24	GeometryPointSize	Geometry
24	Geometry stage exports point size	Geometry
25	ImageGatherExtended	Shader
23	Uses texture gather with non-constant or	Simuli
	independent offsets	
27	StorageImageMultisample	Shader
-,	Uses multi-sample images for non-sampled	~11WW2
	images.	
28	UniformBufferArrayDynamicIndexing	Shader
-	Block-decorated arrays in uniform storage	
	classes use dynamically uniform indexing.	
29	SampledImageArrayDynamicIndexing	Shader
	Arrays of sampled images use dynamically	
	uniform indexing.	
30	StorageBufferArrayDynamicIndexing	Shader
	BufferBlock-decorated arrays in uniform	
	storage classes use dynamically uniform	
	indexing.	
31	StorageImageArrayDynamicIndexing	Shader
	Arrays of non-sampled images are accessed	
1	with dynamically uniform indexing.	
	, – – – – – – – – – – – – – – – – – – –	

	Capability	Depends On
32	ClipDistance	Shader
	Uses the ClipDistance BuiltIn.	
33	CullDistance	Shader
	Uses the CullDistance BuiltIn.	
34	ImageCubeArray	SampledCubeArray
	Uses the Cube Dim with the <i>Arrayed</i>	
	operand in OpTypeImage, without a sampler.	
35	SampleRateShading	Shader
	Uses per-sample rate shading.	
36	ImageRect	SampledRect
	Uses the Rect Dim without a sampler.	
37	SampledRect	Shader
	Uses the Rect Dim with a sampler.	
38	GenericPointer	Addresses
	Uses the Generic Storage Class.	
39	Int8	Kernel
	Uses OpTypeInt to declare 8-bit integer	
	types.	
40	InputAttachment	Shader
	Uses the SubpassData Dim.	
41	SparseResidency	Shader
	Uses OpImageSparse instructions.	
42	MinLod	Shader
	Uses the MinLod Image Operand.	
43	Sampled1D	Shader
	Uses the 1D Dim with a sampler.	
44	Image1D	Sampled1D
	Uses the 1D Dim without a sampler.	
45	SampledCubeArray	Shader
	Uses the Cube Dim with the <i>Arrayed</i>	
	operand in OpTypeImage, with a sampler.	
46	SampledBuffer	Shader
	Uses the Buffer Dim without a sampler.	
47	ImageBuffer	SampledBuffer
	Uses the Buffer Dim without a sampler.	
48	ImageMSArray	Shader
	An MS operand in OpTypeImage indicates	
- 10	multisampled, used without a sampler.	
49	StorageImageExtendedFormats	Shader
	One of a large set of more advanced image	
	formats are used, namely one of those in the	
	Image Format table listed as requiring this	
70	capability.	
50	ImageQuery	Shader
	The sizes, number of samples, or lod, etc. are	
~ .	queried.	GI I
51	DerivativeControl	Shader
	Uses fine or coarse-grained derivatives, e.g.,	
	OpDPdxFine.	

Capability		Depends On
52	InterpolationFunction	Shader
	Uses one of the InterpolateAtCentroid ,	
	InterpolateAtSample, or	
	InterpolateAtOffset GLSL.std.450 extended	
	instructions.	
53	TransformFeedback	Shader
	Uses the Xfb Execution Mode.	
54	GeometryStreams	Geometry
	Uses multiple numbered streams for	
	geometry-stage output.	
55	StorageImageReadWithoutFormat	Shader
	OpImageRead can use the Unknown Image	
	Format for	
56	StorageImageWriteWithoutFormat	Shader
	OpImageWrite can use the Unknown Image	
	Format.	
57	MultiViewport	Geometry
	Multiple viewports are used.	

3.32 Instructions

Form for each instruction:

Opcode Name			Capability
-			Required
Instruction descrip	tion.		Capabilities
1			(when needed)
Word Count is the	high-order 16 bit	ts of word 0 of the	(
	~	Count. If the instruction	
		s, Word Count will also	
	•	inimum size of the	
instruction.	iter stating the in	minimum size of the	
mstruction.			
Oncode is the low	order 16 bits of	word 0 of the	
Opcode is the low-order 16 bits of word 0 of the			
instruction, holding its opcode enumerant.			
Results when present are any Result side or Result Type			
Results, when present, are any Result <id> or Result Type</id>			
created by the instruction. Each one is always 32 bits.			
0 1 1			
Operands, when pr			
instruction's Result			
instruction. Each one is always 32 bits.			
Word Count	Opcode	Results	Operands

3.32.1 Miscellaneous Instructions

OpNop		
This has no semantic impact and can safely be removed from a		
module.		
1	0	

OpUndef			
Make an intermediate object whose value is undefined.			
Result Type is the type of object to make.			
Each consumption of <i>Result <id></id></i> yields an arbitrary, possibly different bit			
pattern.			
3	1	<id>></id>	Result <id></id>
		Result Type	

3.32.2 Debug Instructions

OpSourceContinued

Continue specifying the *Source* text from the previous instruction. This has no semantic impact and can safely be removed from a module.

Continued Source is a continuation of the source text in the previous Source.

The previous instruction must be an OpSource or an OpSourceContinued instruction. As is true for all literal strings, the previous instruction's string was nul terminated. That terminating 0 word from the previous instruction is not part of the source text; the first character of *Continued Source* logically immediately follows the last character of *Source* before its nul.

2 + variable	2	Literal String
		Continued Source

OpSource

Document what source language and text this module was translated from. This has no semantic impact and can safely be removed from a module.

Version is the version of the source language. This literal operand is limited to a single word.

File is an OpString instruction and is the source-level file name.

Source is the text of the source-level file.

Each client API describes what form the *Version* operand takes, per source language.

		1	/ 1	c	
3 + variable	3	Source Language	Literal Number	Optional	Optional
			Version	< <i>id</i> >	Literal String
				File	Source

OpSourceExtension

Document an extension to the source language. This has no semantic impact and can safely be removed from a module.

Extension is a string describing a source-language extension. Its form is dependent on the how the source language describes extensions.

2 + variable	4	Literal String
		Extension

OpName

Assign a name string to another instruction's *Result <id>*. This has no semantic impact and can safely be removed from a module.

Target is the Result $\langle id \rangle$ to assign a name to. It can be the Result $\langle id \rangle$ of any other instruction; a variable, function, type, intermediate result, etc.

Name is the string to assign.

3 + variable	5	<id>></id>	Literal String
		Target	Name

OpMemberName

Assign a name string to a member of a structure type. This has no semantic impact and can safely be removed from a module.

Type is the *<id>* from an OpTypeStruct instruction.

Member is the number of the member to assign in the structure. The first member is member 0, the next is member 1, ... This literal operand is limited to a single word.

Name is the string to assign to the member.

	<u> </u>			
4 + variable	6	< <i>id</i> >	Literal Number	Literal String
		Туре	Member	Name

OpString

Assign a Result < id > to a string for use by other debug instructions (see OpLine and OpSource). This has no semantic impact and can safely be removed from a module. (Removal also requires removal of all instructions referencing Result < id >.)

String is the literal string being assigned a Result <id>.

0	 -			
3 + variable		7	Result <id></id>	Literal String
				String

OpLine

Add source-level location information. This has no semantic impact and can safely be removed from a module.

This location information applies to the instructions physically following this instruction, up to the first occurrence of any of the following: the next end of block, the next **OpLine** instruction, or the next **OpNoLine** instruction.

File must be an OpString instruction and is the source-level file name.

Line is the source-level line number. This literal operand is limited to a single word.

Column is the source-level column number. This literal operand is limited to a single word.

OpLine can generally immediately precede other instructions, with the following exceptions:

- it may not be used until after the annotation instructions, (see the Logical Layout section)
- cannot be the last instruction in a block, which is defined to end with a branch instruction
- if a branch merge instruction is used, the last **OpLine** in the block must be before its merge instruction

		· · · · · · · · · · · · · · · · · · ·		C
4	8	< <i>id</i> >	Literal Number	Literal Number
		File	Line	Column

OpNoLine

Discontinue any source-level location information that might be active from a previous OpLine instruction. This has no semantic impact and can safely be removed from a module.

This instruction can only appear after the annotation instructions (see the Logical Layout section). It cannot be the last instruction in a block, or the second-to-last instruction if the block has a merge instruction. There is not a requirement that there is a preceding **OpLine** instruction.

1	317

3.32.3 Annotation Instructions

OpDecorate					
Add a Decoration to another $\langle id \rangle$.					
_		•	e any <id> that is a forward reference any <id> that is a forward reference and complete any com</id></id>	erence. A set of decorations can be DecorationGroup instruction.	
3 + variable	71	< <i>id</i> >	Decoration	Literal, Literal,	
		Target		See Decoration.	

OpMemberDec	OpMemberDecorate					
Add a Decoration to a member of a structure type.						
Structure type is	Structure type is the <id> of a type from OpTypeStruct.</id>					
Member is the number 1,	<i>Member</i> is the number of the member to decorate in the type. The first member is member 0, the next is member 1					
4 + variable 72 <id> Literal Number Decoration Literal, Literal,</id>						
		Structure Type	Member			
					See Decoration.	

OpDecorationGroup					
A collector for Decorations from OpDecorate instructions. All such OpDecorate instructions targeting this OpDecorationGroup instruction must precede it. Subsequent OpGroupDecorate and OpGroupMemberDecorate instructions consume this instruction's <i>Result <id></id></i> to apply multiple decorations to multiple targets.					
2	73	Result <id></id>			

OpGroupDecorate					
Add a group of Decorations to another < <i>id</i> >.					
Decoration Group is the <id> of an OpDecorationGroup instruction.</id>					
Targets is a list of <id>s to decorate with the groups of decorations.</id>					
2 + variable	74	< <i>id</i> >	<id>, <id>, Targets</id></id>		
		Decoration Group	Targets		

OpGroupMemberDecorate

Add a group of Decorations to members of structure types.

Decoration Group is the *<id>* of an OpDecorationGroup instruction.

Targets is a list of $(\langle id \rangle, Member)$ pairs to decorate with the groups of decorations. Each $\langle id \rangle$ in the pair must be a target structure type, and the associated Member is the number of the member to decorate in the type. The first member is member 0, the next is member 1, ...

2 + variable	75	<id>></id>	<id>, literal,</id>
		Decoration Group	<id>, literal,</id>
			Targets

3.32.4 Extension Instructions

OpExtension

Declare use of an extension to SPIR-V. This allows validation of additional instructions, tokens, semantics, etc.

Name is the extension's name string.

	2 + variable	10	Literal String
			Name

OpExtInstImport

Import an extended set of instructions. It can be later referenced by the *Result <id>*.

Name is the extended instruction-set's name string. There must be an external specification defining the semantics for this extended instruction set.

See Extended Instruction Sets for more information.

3 + variable	11	Result <id></id>	Literal String			
			Name			

OpExtInst

Execute an instruction in an imported set of extended instructions.

Result Type is as defined, per Instruction, in the external specification for Set.

Set is the result of an OpExtInstImport instruction.

Instruction is the enumerant of the instruction to execute within *Set*. This literal operand is limited to a single word. The semantics of the instruction must be defined in the external specification for *Set*.

Operand 1, ... are the operands to the extended instruction.

5 + variable	12	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Literal Number	<id>, <id>,</id></id>
		Result Type		Set	Instruction	
						Operand 1,
						Operand 2,

3.32.5 Mode-Setting Instructions

OpMemor	OpMemoryModel					
Set address	Set addressing model and memory model for the entire module.					
	Addressing Model selects the module's Addressing Model.					
Memory Model selects the module's memory model, see Memory Model.						
3	14	Addressing Model	Memory Model			

OpEntryPoint

Declare an entry point and its execution model.

Execution Model is the execution model for the entry point and its static call tree. See Execution Model.

Entry Point must be the Result <id> of an OpFunction instruction.

Name is a name string for the entry point. A module cannot have two **OpEntryPoint** instructions with the same Execution Model and the same *Name* string.

Interface is a list of <id> of global OpVariable instructions with either Input or Output for its Storage Class operand. These declare the input/output interface of the entry point. They could be a subset of the input/output declarations of the module, and a superset of those referenced by the entry point's static call tree. It is invalid for the entry point's static call tree to reference such an <id> if it was not listed with this instruction.

Interface <*id*> are forward references. They allow declaration of all variables forming an interface for an entry point, whether or not all the variables are actually used by the entry point.

4 + variable	15	Execution Model	<id></id>	Literal String	<id>, <id>,</id></id>
			Entry Point	Name	Interface

OpExecutionMode						
Declare an execution mode for an entry point.						
Entry Point must be the Entry Point <id> operand of an OpEntryPoint instruction. Mode is the execution mode. See Execution Mode.</id>						
3 + variable	3 + variable 16 < id> Execution Mode Optional					
		Entry Point	Mode	literal(s)		
				See Execution		
				Mode		

OpCapability

Declare a capability used by this module.

Capability is the capability declared by this instruction. There are no restrictions on the order in which capabilities are declared.

See the capabilities section for more detail

500	see the capacitates seed on for more detain.				
2	17	7	Capability		
			Capability		

3.32.6 Type-Declaration Instructions

OpTypeVoi	id		
Declare the void type.			
2	19	Result <id></id>	

OpTypeBool

Declare the Boolean type. Values of this type can only be either **true** or **false**. There is no physical size or bit pattern defined for these values. If they are stored (in conjunction with OpVariable), they can only be used with logical addressing operations, not physical, and only with non-externally visible shader Storage Classes: **Workgroup**, **CrossWorkgroup**, **Private**, and **Function**.

Cross workgroup, Titvate, and Function.

- 1			
	2	20	Result <id></id>

OpTypeInt

Declare a new integer type.

Width specifies how many bits wide the type is. This literal operand is limited to a single word. The bit pattern of a signed integer value is two's complement.

Signedness specifies whether there are signed semantics to preserve or validate.

0 indicates unsigned, or no signedness semantics

1 indicates signed semantics.

In all cases, the type of operation of an instruction comes from the instruction's opcode, not the signedness of the operands.

4	21	Result <id></id>	Literal Number	Literal Number	
			Width	Signedness	

OpTypeFloat					
Declare a new floating-point type.					
Width sp	ecifies how	many bits wide the type i	s. The hit pattern of a		
		•	*		
floating-p	point value	is as described by the IEF	EE 754 standard.		
3	22	Result <id></id>	Literal Number		
			Width		

OpTypeVector

Declare a new vector type.

Component Type is the type of each component in the resulting type. It must be a scalar type.

Component Count is the number of components in the resulting type. It must be at least 2.

Components are numbered consecutively, starting with 0.

Components are numbered consecutivery, starting with 0.						
4	23	Result <id></id>	< <i>id</i> >	Literal Number		
			Component Type	Component Count		

OpTypeM	atrix	Capability:		
		Matrix		
Declare a r	new matrix ty	pe.		
Column Ty	pe is the type	e of each column in the	matrix. It must be vector type.	
		1 6 1 1 1		
Column Co	ount is the nu	mber of columns in the	e new matrix type. It must be at least 2.	
Matrix col	ımns are nıın	nhered consecutively s	tarting with 0. This is true	
		•	e	
1 *			he memory layout of a matrix (e.g.,	
KowMajo	r or MatrixS			
4	24	Result <id></id>	< <i>id</i> >	Literal Number
			Column Type	Column Count

OpTypeImage

Declare a new image type. Consumed, for example, by OpTypeSampledImage. This type is opaque: values of this type have no defined physical size or bit pattern.

Sampled Type is the type of the components that result from sampling or reading from this image type. Must be a scalar numerical type or OpTypeVoid.

Dim is the image dimensionality (Dim).

Depth is whether or not this image is a depth image. (Note that whether or not depth comparisons are actually done is a property of the sampling opcode, not of this type declaration.)

- 0 indicates not a depth image
- 1 indicates a depth image
- 2 means no indication as to whether this is a depth or non-depth image

Arrayed must be one of the following indicated values:

- 0 indicates non-arrayed content
- 1 indicates arrayed content

MS must be one of the following indicated values:

- 0 indicates single-sampled content
- 1 indicates multisampled content

Sampled indicates whether or not this image will be accessed in combination with a sampler, and must be one of the following values:

- 0 indicates this is only known at run time, not at compile time
- 1 indicates will be used with sampler
- 2 indicates will be used without a sampler (a storage image)

Image Format is the Image Format, which can be Unknown, depending on the client API.

If Dim is **SubpassData**, *Sampled* must be 2, *Image Format* must be **Unknown**, and the **Execution Model** must be **Fragment**.

Access Qualifier is an image Access Qualifier.

1100000 2.	Tiecess guillyte. Is all mage Tiecess qualities.									
9 +	25	Result	< <i>id</i> >	Dim	Literal	Literal	Literal	Literal	Image	Optional
variable		<id></id>	Sampled		Number	Number	Number	Number	Format	Access
			Туре		Depth	Arrayed	MS	Sampled		Quali-
										fier

OpTypeSampler					
	Declare the sampler type. Consumed by OpSampledImage. This type is opaque: values of this type have no defined physical size or bit pattern.				
2	26	Result <id></id>			

OpTypeSampledImage

Declare a sampled image type, the *Result Type* of OpSampledImage, or an externally combined sampler and image. This type is opaque: values of this type have no defined physical size or bit pattern.

Image Type must be an OpTypeImage. It is the type of the image in the combined sampler and image type.

			•	-	
3	27	Result <id></id>		< <i>id</i> >	
				Image Type	

OpTypeArray

Declare a new array type: a dynamically-indexable ordered aggregate of elements all having the same type.

Element Type is the type of each element in the array.

Length is the number of elements in the array. It must be at least 1. *Length* must come from a constant instruction of an integer-type scalar whose value is at least 1.

Array elements are number consecutively, starting with 0.

	•		•		
4		28	Result <id></id>	< <i>id</i> >	<id></id>
				Element Type	Length

OpTypeRuntin	meArray	Capability:	
Declare a new time.	run-time array ty	Shader	
Element Type is concrete type.	s the type of each	n element in the array. It must be a	
See OpArrayLo	ength for getting		
Objects of this	type can only be		
Uniform Stora	ge Class.		
3	29	Result <id></id>	<id>></id>
			Element Type

OpTypeStruct

Declare a new structure type: an aggregate of potentially heterogeneous members.

Member N type is the type of member N of the structure. The first member is member 0, the next is member $1, \ldots$

If an operand is not yet defined, it must be defined by an OpTypePointer, where the type pointed to is an OpTypeStruct.

2 + variable	30	Result <id></id>	<id>, <id>,</id></id>
			Member 0 type,
			member 1 type,

OpTypeOpaque			Capability:
			Kernel
Declare a structure type with no body			
specified.		-	
3 + variable	31	Result <id></id>	Literal String
			The name of the
			opaque type.

OpTypePointer

Declare a new pointer type.

Storage Class is the Storage Class of the memory holding the object pointed to. If there was a forward reference to this type from an OpTypeForwardPointer, the Storage Class of that instruction must equal the Storage Class of this instruction.

Type is the type of the object pointed to.

4	32	Result <id></id>	Storage Class	<id>></id>
				Туре

OpTypeFunction

Declare a new function type.

OpFunction will use this to declare the return type and parameter types of a function. **OpFunction** is the only valid use of **OpTypeFunction**.

Return Type is the type of the return value of functions of this type. It must be a concrete or abstract type, or a pointer to such a type. If the function has no return value, *Return Type* must be OpTypeVoid.

Parameter N Type is the type $\langle id \rangle$ of the type of parameter N.

Tententere It Type Is	the type there	or the type of purume		
3 + variable	33	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>
			Return Type	Parameter 0 Type,
				Parameter 1 Type,

OpTypeEvent		Capability:
		Kernel
Declare an OpenCL event		
object.		
2	34	Result <id></id>

OpTypeDeviceEvent		Capability:
		DeviceEnqueue
Declare an OpenCL		_
device-side event object.		
2	35	Result <id></id>

OpTypeReserveId		Capability:
		Pipes
Declare an OpenCL		
reservation id object.		
2	36	Result <id></id>

OpTypeQueu	ie	Capability: DeviceEnqueue
Declare an OpenCL queue object.		1
2	37	Result <id></id>

OpTypePipe			Capability:
Declare an OpenCL pipe object type. Qualifier is the pipe access qualifier.			Pipes
3	38	Result <id></id>	Access Qualifier
			Qualifier

OpTypeForwa	rdPointer		Capability:
Declare the Sto	orage Class for	Addresses	
Pointer Type is The type of obj OpTypePointe OpTypeStruct i Storage Class i	ect the pointer er instruction, n nstructions can		
pointed to.			
3	39	< <i>id</i> >	Storage Class
		Pointer Type	

3.32.7 Constant-Creation Instructions

OpConstantTrue				
Declare	Declare a true Boolean-type scalar constant.			
Result T	Result Type must be the scalar Boolean type.			
3	41	< <i>id</i> >	Result <id></id>	
		Result Type		

OpConstantFalse				
Declare a false Boolean-type scalar constant.				
Result Type must be the scalar Boolean type.				
3	42	< <i>id</i> >	Result <id></id>	
		Result Type		

OpConstant

Declare a new integer-type or floating-point-type scalar constant.

Result Type must be a scalar integer type or floating-point type.

Value is the bit pattern for the constant. Types 32 bits wide or smaller take one word. Larger types take multiple words, with low-order words appearing first.

3 + variable 43		< <i>id</i> >	Result <id></id>	Literal, Literal,
		Result Type		Value

OpConstantComposite

Declare a new composite constant.

Result Type must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the *Constituents*. The ordering must be the same between the top-level types in *Result Type* and the *Constituents*.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one *Constituent* for each top-level member/element/component/column of the result. The *Constituents* must appear in the order needed by the definition of the *Result Type*. The *Constituents* must all be <*id*>*s* of other constant declarations.

Ī	3 + variable 44		< <i>id</i> >	Result <id></id>	< <i>id</i> >, < <i>id</i> >,
			Result Type		Constituents

OpCor	nstantSa	mpler			Capability: LiteralSampler	
Declare	Declare a new sampler constant.					
Result	<i>Type</i> mus	et be OpTypeSampler	·.			
Sampler Addressing Mode is the addressing mode; a literal from Sampler Addressing Mode.						
0: Non	is one of Normali malized	•				
Sample	er Filter N	Mode is the filter mod	le; a literal from Sar	npler Filter Mode.		
6	45	<id> Result Type</id>	Result <id></id>	Sampler Addressing Mode	Literal Number Param	Sampler Filter Mode

OpConstantNull

Declare a new null constant value.

The *null* value is type dependent, defined as follows:

- Scalar Boolean: false
- Scalar integer: 0
- Scalar floating point: +0.0 (all bits 0)
- All other scalars: Abstract
- Composites: Members are set recursively to the null constant according to the null value of their constituent types.

Result Type must be one of the following types:

- Scalar or vector Boolean type
- Scalar or vector integer type
- Scalar or vector floating-point type
- Pointer type
- Event type
- Device side event type
- Reservation id type
- Queue type
- Composite type

3	46	< <i>id</i> >	Result <id></id>			
		Result Type				

OpSpecConstantTrue

Declare a Boolean-type scalar specialization constant with a default value of true.

This instruction can be specialized to become either an OpConstantTrue or OpConstantFalse instruction.

Result Type must be the scalar Boolean type.

See Specialization.

_ L				
	3	48	< <i>id</i> >	Result <id></id>
			Result Type	

OpSpecConstantFalse

Declare a Boolean-type scalar specialization constant with a default value of **false**.

This instruction can be specialized to become either an OpConstantTrue or OpConstantFalse instruction.

Result Type must be the scalar Boolean type.

See Specialization.

3	49	< <i>id</i> >	Result <id></id>				
		Result Type					

OpSpecConstant

Declare a new integer-type or floating-point-type scalar specialization constant.

Result Type must be a scalar integer type or floating-point type.

Value is the bit pattern for the default value of the constant. Types 32 bits wide or smaller take one word. Larger types take multiple words, with low-order words appearing first.

This instruction can be specialized to become an OpConstant instruction.

See Specialization.

	*				
3 + variable 50		< <i>id</i> >	Result <id></id>	Literal, Literal,	
			Result Type		Value

OpSpecConstantComposite

Declare a new composite specialization constant.

Result Type must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the *Constituents*. The ordering must be the same between the top-level types in *Result Type* and the *Constituents*.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one *Constituent* for each top-level member/element/component/column of the result. The *Constituents* must appear in the order needed by the definition of the type of the result. The *Constituents* must be the *<id>* of other specialization constant or constant declarations.

This instruction will be specialized to an OpConstantComposite instruction.

See Specialization.

3 + variable	51	< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Constituents

OpSpecConstantOp

Declare a new specialization constant that results from doing an operation.

Result Type must be the type required by the Result Type of Opcode.

Opcode must be one of the following opcodes. This literal operand is limited to a single word.

OpSConvert, **OpFConvert**

OpSNegate, OpNot

OpIAdd, OpISub

OpIMul, OpUDiv, OpSDiv, OpUMod, OpSRem, OpSMod

OpShiftRightLogical, OpShiftRightArithmetic, OpShiftLeftLogical

OpBitwiseOr, OpBitwiseXor, OpBitwiseAnd

OpVectorShuffle, OpCompositeExtract, OpCompositeInsert

OpLogicalOr, OpLogicalAnd, OpLogicalNot,

OpLogicalEqual, OpLogicalNotEqual

OpSelect

OpIEqual

OpULessThan, OpSLessThan

OpUGreaterThan, OpSGreaterThan

OpULess Than Equal, OpSLess Than Equal

OpUGreater Than Equal, OpSGreater Than Equal

If the **Shader** capability was declared, the following opcode is also valid:

OpQuantizeToF16

If the **Kernel** capability was declared, the following opcodes are also valid:

OpConvertFToS, OpConvertSToF

OpConvertFToU, OpConvertUToF

OpUConvert

 $OpConvertPtrToU,\,OpConvertUToPtr$

 $OpGeneric Cast To Ptr, \, OpPtr Cast To Generic \,$

OpBitcast

OpFNegate

OpFAdd, OpFSub

OpFMul, OpFDiv

OpFRem, OpFMod

OpAccessChain, OpInBoundsAccessChain

OpPtrAccess Chain, OpInBoundsPtrAccess Chain

Operands are the operands required by *opcode*, and satisfy the semantics of *opcode*. In addition, all *Operands* must be the < id > s of other constant instructions.

See Specialization.

4 + variable	52	<id></id>	Result <id></id>	Literal Number	<id>, <id>,</id></id>
		Result Type		Opcode	Operands

3.32.8 Memory Instructions

OpVariable

Allocate an object in memory, resulting in a pointer to it, which can be used with OpLoad and OpStore.

Result Type must be an OpTypePointer. Its Type operand is the type of object in memory.

Storage Class is the Storage Class of the memory holding the object. It cannot be Generic.

Initializer is optional. If *Initializer* is present, it will be the initial value of the variable's memory content. *Initializer* must be an <*id*> from a constant instruction. *Initializer* must have the same type as the type pointed to by *Result Type*.

4 + variable	59	< <i>id</i> >	Result <id></id>	Storage Class	Optional
		Result Type			< <i>id</i> >
					Initializer

OpImageTexelPointer

Form a pointer to a texel of an image. Use of such a pointer is limited to atomic operations.

Result Type must be an OpTypePointer whose Storage Class operand is **Image**. Its Type operand must be a scalar numerical type or OpTypeVoid.

Image must be an OpTypePointer with *Type* OpTypeImage. The *Sampled Type* of the type of *Image* must be the same as the *Type* pointed to by *Result Type*. The Dim operand of *Type* cannot be **SubpassData**.

Coordinate and Sample specify which texel and sample within the image to form a pointer to.

Coordinate must be a scalar or vector of integer type. It must have the number of components specified below, given the following *Arrayed* and Dim operands of the type of the OpTypeImage.

If *Arrayed* is 0:

1D: scalar

2D: 2 components 3D: 3 components Cube: 3 components Rect: 2 components Buffer: scalar

If *Arrayed* is 1: **1D**: 2 components **2D**: 3 components **Cube**: 4 components

Sample must be an integer type scalar. It specifies which sample to select at the given coordinate. It must be a valid $\langle id \rangle$ for the value 0 if the OpTypeImage has MS of 0.

6	60	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Image	Coordinate	Sample

OpLoad

Load through a pointer.

Result Type is the type of the loaded object.

Pointer is the pointer to load through. It must be an OpTypePointer whose Type operand is the same as Result Type.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

4 + variable	61	<id></id>	Result <id></id>	<id></id>	Optional
		Result Type		Pointer	Memory Access

OpStore

Store through a pointer.

Pointer is the pointer to store through. It must be an OpTypePointer whose *Type* operand is the same as the type of *Object*.

Object is the object to store.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

3 + variable	62	< <i>id</i> >	< <i>id</i> >	Optional
		Pointer	Object	Memory Access

OpCopyMemory

Copy from the memory pointed to by *Source* to the memory pointed to by *Target*. Both operands must be non-void pointers of the same type. Matching Storage Class is not required. The amount of memory copied is the size of the type pointed to.

Memory Access must be a Memory Access literal. If not present, it is the same as specifying None.

•			1 7 .	
3 + variable	63	< <i>id</i> >	< <i>id</i> >	Optional
		Target	Source	Memory Access

Signedness of 1 and unsigned, and if its	ory point of bytes to stant value to have the value is 0,	copy. It must have cannot be 0. It is a sign bit set. Oth no memory access	e a scalar integer type invalid for both the erwise, as a run-time s will be made.	•	as
4 + variable	64	<id> Target</id>	<id> Source</id>	<id> Size</id>	Optional Memory Access

OpAccessChain

Create a pointer into a composite object that can be used with OpLoad and OpStore.

Result Type must be an OpTypePointer. Its Type operand must be the type reached by walking the Base's type hierarchy down to the last provided index in Indexes, and its Storage Class operand must be the same as the Storage Class of Base.

Base must be a pointer, pointing to the base of a composite object.

Indexes walk the type hierarchy to the desired depth, potentially down to scalar granularity. The first index in *Indexes* will select the top-level member/element/component/element of the base composite. All composite constituents use zero-based numbering, as described by their **OpType...** instruction. The second index will apply similarly to that result, and so on. Once any non-composite type is reached, there must be no remaining (unused) indexes. Each of the *Indexes* must:

- be a scalar integer type,
- be an OpConstant when indexing into a structure.

1					
4 + variable	65	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>
		Result Type		Base	Indexes

OpInBounds	OpInBoundsAccessChain						
Has the same semantics as OpAccessChain, with the addition that the resulting pointer is known to point within the base object.							
4 + variable	66	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>		
		Result Type		Base			
					Indexes		

OpPtrAccessCh	OpPtrAccessChain						
Has the same ser operand.	mantics a	n of the <i>Element</i>	Addresses				
address of the fir computed to be t	Element is used to do the initial dereference of Base: Base is treated as the address of the first element of an array, and the Element element's address is computed to be the base for the Indexes, as per OpAccessChain. The type of Base after being dereferenced with Element is still the same as the original type of Base.						
operation is to se	elect an e		inter an array, and ay, OpAccessChaine array element.				
5 + variable	67	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id>, <id>,</id></id>	
		Result Type		Base	Element		
						Indexes	

OpArr	OpArrayLength							
Length	of a run-tir	Shader						
Result T	Result Type must be an OpTypeInt with 32-bit Width and 0 Signedness.							
Structur	r is a run-time array.							
Array n	<i>nember</i> is tl	ne last member number o	f Structure and must have	e a type from				
OpTypeRuntimeArray.								
5	68	<id></id>	Result <id></id>	< <i>id</i> >	Literal Number			
		Result Type		Structure	Array member			

OpGeneri	cPtrMemSe	mantics		Capability: Kernel
	valid Memor ne specific (n			
Pointer mu	ast point to G	eneric Storage Class.		
Result Type	e must be an			
4	69	<id>></id>	Result <id></id>	<id></id>
		Result Type		Pointer

OpInBoundsPtrAccessChain			Capability: Addresses			
Has the same semantics as OpPtrAccessChain, with the addition that the resulting pointer is known to point within the base object.						
5 + variable	70	<id><id><</id></id>	Result <id></id>	<id> Base</id>	<id> Element</id>	<id>, <id>,</id></id>
		Resuit Type		Dasc	Liemeni	Indexes

3.32.9 Function Instructions

OpFunction

Add a function. This instruction must be immediately followed by one OpFunctionParameter instruction per each formal parameter of this function. This function's body or declaration will terminate with the next OpFunctionEnd instruction.

The *Result <id>* cannot be used generally by other instructions. It can only be used by OpFunctionCall, OpEntryPoint, and decoration instructions.

Result Type must be the same as the Return Type declared in Function Type.

Function Type is the result of an OpTypeFunction, which declares the types of the return value and parameters of the function.

5	54	<id></id>	Result <id></id>	Function Control	< <i>id</i> >
		Result Type			Function Type

OpFunctionParameter

Declare a formal parameter of the current function.

Result Type is the type of the parameter.

This instruction must immediately follow an OpFunction or OpFunctionParameter instruction. The order of contiguous **OpFunctionParameter** instructions is the same order arguments will be listed in an OpFunctionCall instruction to this function. It is also the same order in which *Parameter Type* operands are listed in the OpTypeFunction of the *Function Type* operand for this function's OpFunction instruction.

1 71		×1 1	L	
3	55	< <i>id</i> >		Result <id></id>
		Result Type		

OpFunctionEnd	
Last instruction of a function.	
1	56

OpFunctionCall

Call a function.

Result Type is the type of the return value of the function. It must be the same as the *Return Type* operand of the *Function Type* operand of the *Function* operand.

Function is an OpFunction instruction. This could be a forward reference.

Argument N is the object to copy to parameter N of Function.

Note: A forward call is possible because there is no missing type information: *Result Type* must match the *Return Type* of the function, and the calling argument types must match the formal parameter types.

->F	,	7,1		F	
4 + variable	57	< <i>id</i> >	Result <id></id>	< <i>id</i> >	<id>, <id>,</id></id>
		Result Type		Function	Argument 0,
					Argument 1,

3.32.10 Image Instructions

OpSampledImage

Create a sampled image, containing both a sampler and an image.

Result Type must be the OpTypeSampledImage type.

Image is an object whose type is an OpTypeImage, whose *Sampled* operand is 0 or 1, and whose Dim operand is not **SubpassData**.

Sampler must be an object whose type is OpTypeSampler.

İ	5	86	<id></id>	Result <id></id>	<id></id>	<id></id>
			Result Type		Image	Sampler

OpImageSampl	leImp	licitLod				Capability:	
Sample an image	e with	an implicit leve	el of detail.			Shader	
Result Type must type. Its compon OpTypeImage (u	nents	must be the same	ing				
Sampled Image 1	must l	oe an object who	e.				
Coordinate must array layer]) as than needed, but	neede	a vector larger					
Image Operands	enco	des what operan	ds follow, as per	r Image Operand	ls.		
This instruction consumes an imp		ldition, it					
5 + variable	87	<id> Result Type</id>	Result <id></id>	<id> Sampled</id>	<id><id>Coordinate</id></id>	Optional Image	Optional < <i>id</i> >, < <i>id</i> >,
				Image		Operands	

OpImageSampleExplicitLod

Sample an image using an explicit level of detail.

Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as *Sampled Type* of the underlying OpTypeImage (unless that underlying *Sampled Type* is **OpTypeVoid**).

Sampled Image must be an object whose type is OpTypeSampledImage.

Coordinate must be a scalar or vector of floating-point type or integer type. It contains $(u[, v] ... [, array \, layer])$ as needed by the definition of Sampled Image. Unless the **Kernel** capability is being used, it must be floating point. It may be a vector larger than needed, but all unused components will appear after all used components.

Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present.

7 +	88	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	Image	< <i>id</i> >	Optional
variable		Result	<id></id>	Sampled	Coordinate	Operands		< <i>id</i> >,
		Туре		Image				< <i>id</i> >,

OpImageSa	ampleI	PrefImplicitL	od				Capability: Shader		
Sample an i	mage d	oing depth-co	mparison wit	h an implicit le	vel of detail.		Shauci		
		e a scalar of in		floating-point te.	type. It must be	e the same as			
Sampled Im	<i>age</i> mu								
Coordinate array layer] than needed	or larger								
D_{ref} is the d	epth-co	omparison refe	erence value.						
Image Oper	ands er	ncodes what o	perands follo	w, as per Image	Operands.				
This instruc									
6 +	89	Optional	Optional						
		Result	<id></id>	Sampled	Coordinate	D_{ref}	Image	< <i>id</i> >,	
variable		Type	1	Image		1	Operands	<id>,</id>	

OpImageS	ample	eDrefExplici	tLod					Capability:	
Sample an i Result Type Sampled Type	e same as	Shader							
Sampled Im	age m								
Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v],, [, array layer])$ as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. D_{ref} is the depth-comparison reference value.									
D _{ref} is the d	срш-с	zomparison n	ciciciice varu	ic.					
0 1	<i>Image Operands</i> encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present.								
8 +	90	<id>></id>	Result	<id>></id>	<id>></id>	<id></id>	Image	<id>></id>	Optional
variable		Result Type	<id></id>	Sampled Image	Coordinate	D_{ref}	Operands		< <i>id</i> >, < <i>id</i> >,

OpImageSampleProjImplicitLod Capability: Shader Sample an image with with a project coordinate and an implicit level of detail. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector of four components containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be (u/q [, v/q] [,w/q]), as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 5 + variable 91 <*id*> Result <id> Optional Optional <*id*> <*id*> Result Type Sampled Coordinate <*id*>, <*id*>, Image Image **Operands** . . .

OpImageSampleProjExplicitLod Capability: Shader Sample an image with a project coordinate using an explicit level of detail. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is **OpTypeVoid**). Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector of four components containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be $(u/q \lceil v/q \rceil \lceil w/q \rceil)$, as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. 7 + <*id*> Optional 92 <*id*> Result $\langle id \rangle$ Image $\langle id \rangle$ variable Result < id >Sampled Coordinate Operands <*id*>, Type Image <id>, ...

OpImageSampleProjDrefImplicitLod Capability: Shader Sample an image with a project coordinate, doing depth-comparison, with an implicit level of detail. Result Type must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector of four components containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be $(u/q \lceil v/q \rceil \lceil w/q \rceil)$, as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. D_{ref}/q is the depth-comparison reference value. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. Optional Optional 6+ 93 <*id*> Result <*id*> <*id*> <*id*> variable Result < id >Sampled Coordinate D_{ref} **Image** <*id*>, Operands <*id*>, ... Type *Image*

OpImageSampleProjDrefExplicitLod Capability: Shader Sample an image with a project coordinate, doing depth-comparison, using an explicit level of detail. Result Type must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. The Dim operand of the underlying OpTypeImage must be 1D, 2D, 3D, or Rect, and the Arrayed and MS operands must be 0. Coordinate is a floating-point vector of four components containing (u [, v] [, w], q), as needed by the definition of Sampled Image, with the q component consumed for the projective division. That is, the actual sample coordinate will be $(u/q \lceil v/q \rceil, v/q)$, as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. D_{ref}/q is the depth-comparison reference value. Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. <id> <*id*> Optional 8 + 94 Result <*id*> <*id*> <*id*> Image Result Coordinate Operands <*id*>, variable < id >Sampled D_{ref} Image <*id*>, . . . Type

OpImageFetch

Fetch a single texel from a sampled image.

Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as *Sampled Type* of the underlying OpTypeImage (unless that underlying *Sampled Type* is **OpTypeVoid**).

Image must be an object whose type is OpTypeImage. Its Dim operand cannot be **Cube**, and its *Sampled* operand must be 1.

Coordinate is an integer scalar or vector containing $(u[, v] \dots [, array \, layer])$ as needed by the definition of Sampled Image.

Image Operands encodes what operands follow, as per Image Operands.

0 1			, 1				
5 + variable	95	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Result Type		Image	Coordinate	Image	< <i>id</i> >, < <i>id</i> >,
				_		Operands	
						_	

OpImageGather Capability: Shader Gathers the requested component from four texels. Result Type must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is **OpTypeVoid**). It has one component per gathered texel. Sampled Image must be an object whose type is OpTypeSampledImage. Its OpTypeImage must have a Dim of 2D, Cube, or Rect. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [,$ array layer]) as needed by the definition of Sampled Image. Component is the component number that will be gathered from all four texels. It must be 0, 1, 2 or 3. Image Operands encodes what operands follow, as per Image Operands. Optional <*id*> Result <*id*> <*id*> <*id*> Optional 6+ 96 variable Result < id >Sampled Coordinate Component **Image** <*id*>, Type Image Operands <*id*>, ...

OpImageDr	efGat	her					Capability:	
Gathers the r	equest	ed depth-comp	parison from fo	our texels.			Shader	
Result Type n Sampled Type								
Sampled Ima OpTypeImag	_							
		e a scalar or ve			t contains (u[,	v] [,		
		mparison refer		as per Image	Operands			
6+	<id>></id>	Optional	Optional					
variable								< <i>id</i> >,
		Type Image						<id>,</id>

OpImageRead

Read a texel from an image without a sampler.

Result Type must be a scalar or vector of floating-point type or integer type. Its component type must be the same as Sampled Type of the OpTypeImage (unless that Sampled Type is OpTypeVoid).

Image must be an object whose type is OpTypeImage with a *Sampled* operand of 0 or 2. If the *Sampled* operand is 2, then some dimensions require a capability; e.g., one of **Image1D**, **ImageRect**, **ImageBuffer**, **ImageCubeArray**, or **ImageMSArray**.

Coordinate is an integer scalar or vector containing non-normalized texel coordinates ($u[, v] \dots [, array \, layer]$) as needed by the definition of Image. If the coordinates are outside the image, the memory location that is accessed is undefined.

When the *Image* Dim operand is **SubpassData**, *Coordinate* is relative to the current fragment location. That is, the integer value (rounded down) of the current fragment's window-relative (x, y) coordinate is added to (u, v).

When the *Image* Dim operand is not **SubpassData**, the **Image** Format must not be **Unknown**, unless the **StorageImageReadWithoutFormat** Capability was declared.

Image Operands encodes what operands follow, as per Image Operands.

5 + variable	98	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Result Type		Image	Coordinate	Image	<id>, <id>,</id></id>
						Operands	
						_	

OpImageWrite

Write a texel to an image without a sampler.

Image must be an object whose type is OpTypeImage with a *Sampled* operand of 0 or 2. If the *Sampled* operand is 2, then some dimensions require a capability; e.g., one of **Image1D**, **ImageRect**, **ImageBuffer**, **ImageCubeArray**, or **ImageMSArray**. Its Dim operand cannot be **SubpassData**.

Coordinate is an integer scalar or vector containing non-normalized texel coordinates ($u[, v] \dots [, array \, layer]$) as needed by the definition of Image. If the coordinates are outside the image, the memory location that is accessed is undefined.

Texel is the data to write. Its component type must be the same as *Sampled Type* of the OpTypeImage (unless that *Sampled Type* is **OpTypeVoid**).

The Image Format must not be Unknown, unless the StorageImageWriteWithoutFormat Capability was declared.

Image Operands encodes what operands follow, as per Image Operands.

1 THE SC OF CHILL		, what operands for	, as per	operanos.		
4 + variable	99	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	Optional	Optional
		Image	Coordinate	Texel	Image	<id>, <id>,</id></id>
					Operands	

OpImag	OpImage								
Extract	Extract the image from a sampled image.								
Result T	<i>ype</i> must	be OpTypeImage.							
_ ^	0	ust have type OpTypeS	ampledImage whose Im	nage Type is the same					
as <i>Result Type</i> . 4 100 <id> Result <id> <id> </id></id></id>									
7	100	Result Type	Result \id>	Sampled Image					

OpImag	geQueryFor	mat		Capability:
Query th	ne image form	Kernel		
	<i>Type</i> must be a Channel Data	m		
Image n	nust be an obj			
4	101	<id>></id>		
		Result Type		Image

OpImag	eQueryOrde	r		Capability:
Result Ty	pe must be a	scalar integer type. The	vith an Unknown Image Format . resulting value is an enumerant from	Kernel
		ect whose type is OpTyp		
4	102	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Image

OpImag	geQueryS	izeLod			Capability:
Query th	ne dimensi	ions of <i>Image</i> for mipm	nap level for <i>Level of D</i>	etail.	Kernel, ImageQuery
Result T	<i>ype</i> must l	be an integer type scala	r or vector. The number	r of components must	be
1 for 1D	Dim,				
2 for 2D	, and Cub	e Dimensionalities,			
3 for 3D	Dim,				
plus 1 m	nore if the	image type is arrayed.	This vector is filled in	with (width [, height] [,
depth] [,	, elements) where <i>elements</i> is the	e number of layers in a	n image array, or the nu	umber
of cubes	in a cube	-map array.			
2D, 3D,		and its MS must be 0.	pTypeImage. Its Dim o See OpImageQuerySize		
Level of specifica		used to compute which	mipmap level to query	, as described in the Al	PI
5	103	<id>></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Image	Level of Detail

OpImageQueryS	ize		Capability:
Query the dimensi	ons of <i>Image</i> , with no level	l of detail.	Kernel, ImageQuery
Result Type must l	be an integer type scalar or	vector. The number of components i	must
be			
1 for Buffer Dim ,			
2 for 2D and Rect	Dimensionalities,		
3 for 3D Dim,			
plus 1 more if the	image type is arrayed. This	s vector is filled in with (width [, height	ght]
[, elements]) wher	e elements is the number of	f layers in an image array.	
		peImage. Its Dim operand must be o	
of Rect or Buffer ,	or if its MS is 1, it can be 2	2D , or, if its <i>Sampled Type</i> is 0 or 2,	it
can be 2D or 3D .	t cannot be an image with	level of detail; there is no implicit	
level-of-detail con	sumed by this instruction. S	See OpImageQuerySizeLod for	
querying images h	aving level of detail.		
4 104	<id></id>	Result <id></id>	< <i>id</i> >
	Result Type		Image

OpImage	QueryLo	od			Capability:
					ImageQuery
Query the	mipmap	level and the level of det	ail for a hypothetica	l sampling of Image at	
Coordinat	te using a	in implicit level of detail.			
Result Typ	e must b	e a two-component floati	ng-point type vector	:	
		nt of the result will contain	U 1 11		
		nent of the result will con		•	he
base level.	-		*		
_		object whose type is OpT	ypeImage. Its Dim o	perand must be one of	1D,
2D , 3D , or	r Cube .				
Coordinat	te must h	e a scalar or vector of flo	ating-point type or i	nteger type. It contains	(uf
		r]) as needed by the defin			. –
		used, it must be floating	•	age. Omess the Herme	-
	8	6	L		
If called o	n an inco	omplete image, the results	s are undefined.		
		only valid in the Fragme		. In addition, it consun	nes an
	erivative	that can be affected by co			
5	105	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Image	Coordinate

OpImage(QueryLevels			Capability:
Result Type	number of mines a second to the must be a second to the API sp	Kernel, ImageQuery		
	t be an objec 3D, or Cube	t whose type is OpTypeImage.	Its Dim operand must be one	
4	106	<id>></id>	Result <id></id>	<id>></id>
		Result Type		Image

OpImageQuerySampl	es		Capability:
Result Type must be a s	imples available per texel fetch calar integer type. The result is it whose type is OpTypeImage.	the number of samples.	Kernel, ImageQuery
4 107	< <i>id</i> >	Result <id></id>	<id>></id>
	Result Type		Image

OpImageSparseSampleImplicitLod Capability: **SparseResidency** Sample a sparse image with an implicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled *Type* is **OpTypeVoid**). Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [,$ array layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. Image Operands encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 5 + variable 305 <*id*> Result <id> <*id*> $\langle id \rangle$ Optional Optional Result Type Sampled Coordinate **Image** <*id*>, <*id*>, Image Operands

OpImageSparseSampleExplicitLod Capability: **SparseResidency** Sample a sparse image using an explicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled Type is OpTypeVoid). Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type or integer type. It contains $(u[,v]...[, array \, layer])$ as needed by the definition of Sampled Image. Unless the Kernel capability is being used, it must be floating point. It may be a vector larger than needed, but all unused components will appear after all used components. *Image Operands* encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. Optional 7+ 306 <*id*> Result $\langle id \rangle$ $\langle id \rangle$ Image $\langle id \rangle$ variable Result < id >Sampled Coordinate Operands <*id*>, Type Image <id>, ...

OpImageSparseSampleDrefImplicitLod Capability: **SparseResidency** Sample a sparse image doing depth-comparison with an implicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains (u[, v], ..., [, v])array layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. D_{ref} is the depth-comparison reference value. *Image Operands* encodes what operands follow, as per Image Operands. This instruction is only valid in the Fragment Execution Model. In addition, it consumes an implicit derivative that can be affected by code motion. 307 <*id*> Result <id> Optional Optional 6+ <*id*> <*id*> variable Result <id> Sampled Coordinate <*id*>, D_{ref} Image <*id*>, ... Type Image **Operands**

OpImageSparseSampleDrefExplicitLod Capability: **SparseResidency** Sample a sparse image doing depth-comparison using an explicit level of detail. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. Sampled Image must be an object whose type is OpTypeSampledImage. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [, array])$ layer]) as needed by the definition of Sampled Image. It may be a vector larger than needed, but all unused components will appear after all used components. D_{ref} is the depth-comparison reference value. Image Operands encodes what operands follow, as per Image Operands. At least one operand setting the level of detail must be present. <*id*> <*id*> <*id*> 8 + 308 Result $\overline{\langle id \rangle}$ Image <*id*> Optional variable Result < id >Sampled Coordinate D_{ref} Operands <*id*>, Type *Image* <*id*>, ...

OpImageSpa	rseSan	npleProjImplio	citLod			Capability:	
				uction is invalid		SparseResid	ency
5 + variable	309	<id> Result Type</id>	Result <id></id>	<id><id><id>Image</id></id></id>	<id> Coordinate</id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,

		ampleProjEx	•	instruction is in	nvalid.		Capability: SparseResid	dency
Sample a spa	arse im	age with a pro	ojective coordi	nate using an e	explicit level of	detail.		
7 +	310	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	Image	< <i>id</i> >	Optional
variable	variable Result <id> Sampled Coordinate Operands </id>							
		Type		Image		_		< <i>id</i> >,

OpImageS	SparseSa	ampleProjD	PrefImplicitLo	od			Capability: SparseResi	dency
Instruction	reserved	d for future u	use. Use of thi	s instruction is	invalid.			•
Sample a simplicit lev	-		orojective coor	dinate, doing de	epth-comparisor	n, with an		
6 + variable	311	<id> Result Type</id>	Result <id></id>	<id><id><id>Image</id></id></id>	<id>Coordinate</id>	<id> D_{ref}</id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,

OpImage	Sparse	SamplePro	jDrefExplic	itLod				Capability SparseRe	
	sparse i	mage with		f this instructi		nparison, ı	using an		
8 + variable	312	<id> Result Type</id>	Result <id></id>	<id> Sampled Image</id>	<id> Coordinate</id>	$<$ i $d>$ D_{ref}	Image Operands	<id>></id>	Optional < <i>id</i> >, < <i>id</i> >,

OpImageSparseFetch Capability: **SparseResidency** Fetch a single texel from a sparse image. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a vector of four components of floating-point type or integer type. Its components must be the same as Sampled Type of the underlying OpTypeImage (unless that underlying Sampled *Type* is **OpTypeVoid**). Image must be an object whose type is OpTypeImage. Its Dim operand cannot be Cube. Coordinate is an integer scalar or vector containing (u[, v] ... [, array layer]) as needed by the definition of Sampled Image. Image Operands encodes what operands follow, as per Image Operands. 5 + variable 313 <*id*> Result <id> <*id*> Optional Optional <*id*> Result Type Image Coordinate **Image** <*id*>, <*id*>, Operands . . .

Result Type I be an integer OpImageSpa components Sampled Typ OpTypeVoid	request must be type s urseTex of floar e of the d). It have	e an OpType calar. It will elsResident. ting-point type e underlying as one composit be an obje	Struct with tw hold a Reside The second more or integer to OpTypeImagement per gath	is OpTypeSam	e first member' an be passed to a vector of founents must be the inderlying Samp	r ne same as	Capability: SparseResi	dency
Coordinate 1	nust be	a scalar or v		ng-point type.	It contains (u[,	v] [,		
	s the c	•		<i>ampled Image</i> . I be gathered fr	om all four texe	els. It must		
				w, as per Image		. • 1.		0 .: 1
6 + variable	314	<id> Result Type</id>	Result <id></id>	<id><id><id><id>Image</id></id></id></id>	<id><id>Coordinate</id></id>	<id> Component</id>	Optional Image Operands	Optional < <i>id</i> >, < <i>id</i> >,

OpImageSparseDrefGather Capability: **SparseResidency** Gathers the requested depth-comparison from four texels of a sparse image. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar of integer type or floating-point type. It must be the same as Sampled Type of the underlying OpTypeImage. It has one component per gathered texel. Sampled Image must be an object whose type is OpTypeSampledImage. Its OpTypeImage must have a Dim of 2D, Cube, or Rect. Coordinate must be a scalar or vector of floating-point type. It contains $(u[, v] \dots [,$ array layer]) as needed by the definition of Sampled Image. D_{ref} is the depth-comparison reference value. *Image Operands* encodes what operands follow, as per Image Operands. Optional 6+ 315 <*id*> Result $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ $\overline{\langle id \rangle}$ Optional variable Result < id >Sampled Coordinate D_{ref} Image <*id*>, Туре Image Operands <*id*>, ...

OpImageS	SparseTexels	Resident		Capability:
uncommitt	Translates a <i>Resident Code</i> into a Boolean. Result is false if any of the texels were in uncommitted texture memory, and true otherwise. *Result Type must be a Boolean type scalar.			SparseResidency
Resident C	<i>ode</i> is a valudde.			
4	316	<id>></id>	Result <id></id>	<id>></id>
		Result Type		Resident Code

OpImageSparseRead Capability: **SparseResidency** Read a texel from a sparse image without a sampler. Result Type must be an OpTypeStruct with two members. The first member's type must be an integer type scalar. It will hold a Residency Code that can be passed to OpImageSparseTexelsResident. The second member must be a scalar or vector of floating-point type or integer type. Its component type must be the same as Sampled *Type* of the OpTypeImage (unless that *Sampled Type* is **OpTypeVoid**). *Image* must be an object whose type is OpTypeImage with a *Sampled* operand of 2. Coordinate is an integer scalar or vector containing non-normalized texel coordinates $(u[,v]...[, array \, layer])$ as needed by the definition of *Image*. If the coordinates are outside the image, the memory location that is accessed is undefined. The Image Format must not be Unknown, unless the StorageImageReadWithoutFormat Capability was declared. *Image Operands* encodes what operands follow, as per Image Operands. 5 + variable 320 <*id*> Result <id> <*id*> $\langle id \rangle$ Optional Optional Result Type Image Coordinate **Image** <*id*>, <*id*>, Operands

3.32.11 Conversion Instructions

OpConvertFToU

Convert (value preserving) from floating point to unsigned integer, with round toward 0.0.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

Float Value must be a scalar or vector of floating-point type. It must have the same number of components as *Result Type*.

Results are computed per component.

Ī	4	109	< <i>id</i> >	Result <id></id>	<id></id>
			Result Type		Float Value

OpConvertFToS

Convert (value preserving) from floating point to signed integer, with round toward 0.0.

Result Type must be a scalar or vector of integer type.

Float Value must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

Results are computed per component.

4	110	1.	D 10 .2.16	1.
4	110	<1d>	Result <10>	< <i>id></i>
		Result Type		Float Value

OpConvertSToF

Convert (value preserving) from signed integer to floating point.

Result Type must be a scalar or vector of floating-point type.

Signed Value must be a scalar or vector of integer type. It must have the same number of components as Result Type.

	The state of the s			
4	111	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Signed Value

OpConvertUToF

Convert (value preserving) from unsigned integer to floating point.

Result Type must be a scalar or vector of floating-point type.

Unsigned Value must be a scalar or vector of integer type. It must have the same number of components as *Result Type*.

Results are computed per component.

ı	4	112	<id></id>	Result <id></id>	<id></id>
			Result Type		Unsigned Value

OpUConvert

Convert (value preserving) unsigned width. This is either a truncate or a zero extend.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

Unsigned Value must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width cannot equal the component width in *Result Type*.

Results are computed per component.

4	113	< <i>id</i> >	Result <id></id>	< <i>id</i> >	
		Result Type		Unsigned Value	

OpSConvert

Convert (value preserving) signed width. This is either a truncate or a sign extend.

Result Type must be a scalar or vector of integer type.

Signed Value must be a scalar or vector of integer type. It must have the same number of components as Result Type. The component width cannot equal the component width in Result Type.

Results are computed per component.

4	114	< <i>id></i>	Result <id></id>	< <i>id</i> >
		Result Type		Signed Value

OpFConvert

Convert (value preserving) floating-point width.

Result Type must be a scalar or vector of floating-point type.

Float Value must be a scalar or vector of floating-point type. It must have the same number of components as Result Type. The component width cannot equal the component width in Result Type.

	4	115	< <i>id</i> >	Result <id></id>	<id></id>	
			Result Type		Float Value	

OpQuan	tizeToF16			Capability:
Quantize	a floating-poi	Shader e.		
Result Ty	<i>pe</i> must be a s 32 bits.			
Value is t	he value to qu	antize. The type of V_{α}	alue must be the same as Result Type.	
NaN, but large to r Value is t value, the	s an infinity, the not necessarile epresent as a legative with a ceresult is negative as a normalizer			
The Rela	xedPrecision			
Results a	re computed p			
4	116	<id>></id>	Result <id></id>	<id></id>
		Result Type		Value

OpConver	tPtrToU			Capability:
		Addresses		
Convert a pointer to an unsigned integer type. A <i>Result Type</i> width larger than the				
width of Pa	width of <i>Pointer</i> will zero extend. A <i>Result Type</i> smaller than the width of <i>Pointer</i>			
will truncat	te. For same-	width source and result, this is t	he same as OpBitCast.	
			-	
Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.				
4	117	<id>></id>	Result <id></id>	<id></id>
		Result Type		Pointer

Capability:
Kernel
llue of
numbon
number
<id>></id>
Signed Value

OpSatConvertUToS			Capability:	
Convert an unsigned in	nteger to signed integer. Converte	ed values outside the		
representable range of	Result Type are clamped to the n	earest representable value of		
Result Type.				
Result Type must be a	scalar or vector of integer type.			
Unsigned Value must	be a scalar or vector of integer ty	pe. It must have the same		
number of component	s as Result Type.			
Results are computed per component.				
4 119	<id>></id>	Result <id></id>	<id></id>	
	Result Type		Unsigned Value	

Value point Value point Result Type same as Op	integer to po ther will truncater will zero et et must be an oblitCast.	Capability: Addresses		
4	120	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		Integer Value

OpPtrCast	ToGeneric	Capability:		
Convert a po	ointer's Stor	Kernel		
Result Type	must be an	OpTypePointer. Its Storage Cla	ss must be Generic.	
Pointer must	t point to the	e Workgroup, CrossWorkgro	up, or Function Storage Class.	
Result Type	and <i>Pointer</i>			
4	121	< <i>id</i> >	< <i>id</i> >	
		Result Type		Pointer

OpGenerio	CastToPtr	Capability:		
Convert a p	ointer's Stor	Kernel		
1	must be an kgroup , or F			
Pointer mu	st point to th			
Result Type	and <i>Pointer</i>			
4	122	<id>></id>	Result <id></id>	<id>></id>
		Result Type		Pointer

OpGene	Capability:				
	Kernel				
Attempts	s to explici	tly convert Pointer to Sa	torage storage-class po	ointer value.	
Result Ty	<i>pe</i> must b	e an OpTypePointer. Its	Storage Class must be	e Storage.	
		0.50	er tat at	T CD L	
		OpTypePointer whose	* *	* *	
Type.Poir	nter must	point to the Generic St o	orage Class. If the cast	fails, the instruction result	
is an Opt	ConstantN	ull pointer in the Storag	e Storage Class.		
Storage 1	must be on	e of the following litera	l values from Storage	Class: Workgroup,	
CrossWo					
5	123	<id></id>	Result <id></id>	< <i>id</i> >	Storage Class
		Result Type		Pointer	Storage

OpBitcast

Bit pattern-preserving type conversion.

Result Type must be an OpTypePointer, or a scalar or vector of numerical-type.

Operand must be an OpTypePointer, or a scalar or vector of numerical-type. It must be a different type than *Result Type*.

If *Result Type* is a pointer, *Operand* must be a pointer or integer scalar. If *Operand* is a pointer, *Result Type* must be a pointer or integer scalar.

If *Result Type* has the same number of components as *Operand*, they must also have the same component width, and results are computed per component.

If *Result Type* has a different number of components than *Operand*, the total number of bits in *Result Type* must equal the total number of bits in *Operand*. Let L be the type, either *Result Type* or *Operand's* type, that has the larger number of components. Let S be the other type, with the smaller number of components. The number of components in L must be an integer multiple of the number of components in S. The first component (that is, the only or lowest-numbered component) of S maps to the first components of L, and so on, up to the last component of S mapping to the last components of L. Within this mapping, any single component of S (mapping to multiple components of L) maps its lower-ordered bits to the lower-numbered components of L.

4	124	< <i>id</i> >	Result <id></id>	<id></id>
		Result Type		Operand

3.32.12 Composite Instructions

OpVectorExtractDynamic

Extract a single, dynamically selected, component of a vector.

Result Type must be a scalar type.

Vector must be an OpTypeVector whose *Component Type* is *Result Type*.

Index must be a scalar integer 0-based index of which component of *Vector* to extract.

The value read is undefined if *Index's* value is less than zero or greater than or equal to the number of components in *Vector*.

5	77	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Vector	Index

OpVectorInsertDynamic

Make a copy of a vector, with a single, variably selected, component modified.

Result Type must be an OpTypeVector.

Vector must have the same type as Result Type and is the vector that the non-written components will be copied from.

Component is the value that will be supplied for the component selected by *Index*. It must have the same type as the type of components in *Result Type*.

Index must be a scalar integer 0-based index of which component to modify.

What is written is undefined if *Index's* value is less than zero or greater than or equal to the number of components in *Vector*.

6	78	<id></id>	Result <id></id>	<id></id>	< <i>id</i> >	<id></id>
		Result Type		Vector	Component	Index

OpVectorShuffle

Select arbitrary components from two vectors to make a new vector.

Result Type must be an OpTypeVector. The number of components in *Result Type* must be the same as the number of *Component* operands.

Vector 1 and Vector 2 must both have vector types, with the same Component Type as Result Type. They do not have to have the same number of components as Result Type or with each other. They are logically concatenated, forming a single vector with Vector 1's components appearing before Vector 2's. The components of this logical vector are logically numbered with a single consecutive set of numbers from 0 to N - 1, where N is the total number of components.

Components are these logical numbers (see above), selecting which of the logically numbered components form the result. They can select the components in any order and can repeat components. The first component of the result is selected by the first Component operand, the second component of the result is selected by the second Component operand, etc. A Component literal may also be FFFFFFFF, which means the corresponding result component has no source and is undefined. All Component literals must either be FFFFFFFF or in [0, N - 1] (inclusive).

Note: A vector "swizzle" can be done by using the vector for both *Vector* operands, or using an OpUndef for one of the *Vector* operands.

1						
5 + variable	79	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Literal, Literal,
		Result Type		Vector 1	Vector 2	
						Components

OpCompositeConstruct

Construct a new composite object from a set of constituent objects that will fully form it.

Result Type must be a composite type, whose top-level members/elements/components/columns have the same type as the types of the operands, with one exception. The exception is that for constructing a vector, the operands may also be vectors with the same component type as the Result Type component type. When constructing a vector, the total number of components in all the operands must equal the number of components in Result Type.

Constituents will become members of a structure, or elements of an array, or components of a vector, or columns of a matrix. There must be exactly one *Constituent* for each top-level member/element/component/column of the result, with one exception. The exception is that for constructing a vector, a contiguous subset of the scalars consumed can be represented by a vector operand instead. The *Constituents* must appear in the order needed by the definition of the type of the result. When constructing a vector, there must be at least two *Constituent* operands.

3 + variable	80	< <i>id</i> >	Result <id></id>	<id>, <id>,</id></id>
		Result Type		Constituents

OpCompositeExtract

Extract a part of a composite object.

Result Type must be the type of object selected by the last provided index. The instruction result is the extracted object.

Composite is the composite to extract from.

Indexes walk the type hierarchy, potentially down to component granularity, to select the part to extract. All indexes must be in bounds. All composite constituents use zero-based numbering, as described by their **OpType...** instruction.

4 + variable	81	<id>></id>	Result <id></id>	< <i>id</i> >	Literal, Literal,
		Result Type		Composite	Indexes

OpCompositeInsert

Make a copy of a composite object, while modifying one part of it.

Result Type must be the same type as Composite.

Object is the object to use as the modified part.

Composite is the composite to copy all but the modified part from.

Indexes walk the type hierarchy of *Composite* to the desired depth, potentially down to component granularity, to select the part to modify. All indexes must be in bounds. All composite constituents use zero-based numbering, as described by their **OpType...** instruction. The type of the part selected to modify must match the type of *Object*.

5 + variable	82	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	Literal, Literal,
		Result Type		Object	Composite	
						Indexes

OpCopyObject							
Make a copy of <i>Operand</i> . There are no dereferences involved.							
Result Type must match Operand type. There are no other							
restrictions on the types.							
4	83	< <i>id</i> >	Result <id></id>	< <i>id</i> >			
		Result Type		Operand			

OpTransp	ose			Capability:
Transpose	a matrix.	Matrix		
	e must be an ereverse of tho			
Matrix mu	st be an objec	et of type from an OpTypeMatri	x instruction.	
4	84	<id></id>	Result <id></id>	<id></id>
		Result Type		Matrix

3.32.13 Arithmetic Instructions

OpSNegate

Signed-integer subtract of Operand from zero.

Result Type must be a scalar or vector of integer type.

Operand's type must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width must equal the component width in *Result Type*.

Results are computed per component.

İ	4	126	<id></id>	Result <id></id>	<id></id>
			Result Type		Operand

OpFNegate

Floating-point subtract of Operand from zero.

Result Type must be a scalar or vector of floating-point type.

The type of *Operand* must be the same as *Result Type*.

Results are computed per component.

4	127	<id>></id>	Result <id></id>	< <i>id</i> >
		Result Type		Operand

OpIAdd

Integer addition of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

	1	1 1			
5	128	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpFAdd

Floating-point addition of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results	are compi	ned per component.				
5	129	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpISub

Integer subtraction of Operand 2 from Operand 1.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

5	130	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpFSub

Floating-point subtraction of *Operand 2* from *Operand 1*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component.

5	131	< <i>id</i> >	Result <id></id>	<id></id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpIMul

Integer multiplication of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component.

5	132	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpFMul

Floating-point multiplication of *Operand 1* and *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

results	Acoustic are compared per component.					
5	133	< <i>id</i> >	Result <id></id>	< <i>id></i>	< <i>id></i>	
		Result Type		Operand 1	Operand 2	

OpUDiv

Unsigned-integer division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	134	<id></id>	Result <id></id>	<id></id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpSDiv

Signed-integer division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	135	<id></id>	Result <id></id>	<id>></id>	<id>></id>	
		Result Type		Operand 1	Operand 2	

OpFDiv

Floating-point division of *Operand 1* divided by *Operand 2*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	136	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpUMod

Unsigned modulo operation of *Operand 1* modulo *Operand 2*.

Result Type must be a scalar or vector of integer type, whose Signedness operand is 0.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	137	<id>></id>	Result <id></id>	<id>></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpSRem

Signed remainder operation of *Operand 1* divided by *Operand 2*. The sign of a non-0 result comes from *Operand 1*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	138	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpSMod

Signed modulo operation of *Operand 1* modulo *Operand 2*. The sign of a non-0 result comes from *Operand 2*.

Result Type must be a scalar or vector of integer type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	139	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpFRem

Floating-point remainder operation of *Operand 1* divided by *Operand 2*. The sign of a non-0 result comes from *Operand 1*.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	140	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpFMod

Floating-point modulo operation of Operand 1 modulo Operand 2. The sign of a non-0 result comes from Operand 2.

Result Type must be a scalar or vector of floating-point type.

The types of *Operand 1* and *Operand 2* both must be the same as *Result Type*.

Results are computed per component. The resulting value is undefined if *Operand 2* is 0.

5	141	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

Scalar

Op Vector TimesScalar Scale a floating-point vector. Result Type must be a vector of floating-point type. The type of Vector must be the same as Result Type. Each component of Vector is multiplied by Scalar. Scalar must have the same type as the Component Type in Result Type. 5 142 <id><</td> <id><</td>

Vector

Result Type

OpMatr	OpMatrixTimesScalar					
	Scale a floating-point matrix. Result Type must be an OpTypeMatrix whose Column Type is a vector of floating-point type.					
Matrix is	s multiplie	must be the same as <i>Result</i> d by <i>Scalar</i> . the same type as the <i>Comp</i> .				
5						
		Result Type		Matrix	Scalar	

OpVecto	OpVectorTimesMatrix						
		Matrix					
Linear-a	lgebraic Ve	ector X Matrix.					
Result Ty	<i>ype</i> must b						
Vector m	nust be a ve	ector with the same Comp	onent Type as the C	omponent Type in Resu	lt		
Type. Its	number of	f components must equal	the number of comp	onents in each column	in		
Matrix.							
Matrix n	nust be a m	ult					
Type. Its number of columns must equal the number of components in Result Type.							
5	< <i>id</i> >						
		Result Type		Vector	Matrix		

OpMatr	Capability:					
Linear-al	Linear-algebraic <i>Vector X Matrix</i> .					
Result Ty	<i>pe</i> must b	e a vector of floating-poin	nt type.			
Matrix n	nust be an	OpTypeMatrix whose Co	lumn Type is Result Type	e.		
Vector m	nust be a ve	ector with the same Comp	ponent Type as the Comp	onent Type in Result		
Type. Its	Type. Its number of components must equal the number of columns in Matrix.					
5	<id></id>					
		Result Type		Matrix	Vector	

OpMatr	OpMatrixTimesMatrix					
Linear-a	Matrix					
Result Ty	<i>vpe</i> must b	e an OpTypeMatrix whos	e Column Type is a vector	or of floating-point type.		
LeftMatr Type.	LeftMatrix must be a matrix whose Column Type is the same as the Column Type in Result Type.					
RightMatrix must be a matrix with the same Component Type as the Component Type in Result Type. Its number of columns must equal the number of columns in Result Type. Its columns must have the same number of components as the number of columns in LeftMatrix.						
5	5 146 <id> Result <id> <id> </id></id></id>					
		Result Type		LeftMatrix	RightMatrix	

OpOu	OpOuterProduct						
		Matrix					
Linear-							
Result	Result Type must be an OpTypeMatrix whose Column Type is a vector of floating-point type.						
Vector	1 must hav	ve the same type as the	Column Type in Result	Туре.			
Vector	2 must be	a vector with the same	Component Type as the	Component Type in Result			
Type. I	<i>Type</i> . Its number of components must equal the number of columns in <i>Result Type</i> .						
5	<id>></id>						
		Result Type		Vector 1	Vector 2		

OpDo	OpDot								
Dot pi	Dot product of <i>Vector 1</i> and <i>Vector 2</i> .								
Result	<i>Type</i> m	ust be a floating-p	oint type scalar.						
	Vector 1 and Vector 2 must have the same type, and their component type must be Result Type.								
5	5 148 <id> Result <id> <id> <id> </id></id></id></id>								
		Result Type		Vector 1	Vector 2				

OpIAddCarry

Result is the unsigned integer addition of *Operand 1* and *Operand 2*, including its carry.

Result Type must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits (full component width) of the addition.

Member 1 of the result gets the high-order (carry) bit of the result of the addition. That is, it gets the value 1 if the addition overflowed the component width, and 0 otherwise.

5	149	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
	117	Result Type		Operand 1	Operand 2

OpISubBorrow

Result is the unsigned integer subtraction of *Operand 2* from *Operand 1*, and what it needed to borrow.

Result Type must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits (full component width) of the subtraction. That is, if *Operand 1* is larger than *Operand 2*, member 0 gets the full value of the subtraction; if *Operand 2* is larger than *Operand 1*, member 0 gets $2^w + Operand 1 - Operand 2$, where w is the component width.

Member 1 of the result gets 0 if *Operand* 1 > Operand 2, and gets 1 otherwise.

5 150		<id>></id>	Result <id></id>	< <i>id></i>		
		Result Type		Operand 1	Operand 2	

OpUMulExtended

Result is the full value of the unsigned integer multiplication of *Operand 1* and *Operand 2*.

Result Type must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type, whose *Signedness* operand is 0.

Operand 1 and Operand 2 must have the same type as the members of Result Type. These are consumed as unsigned integers.

Results are computed per component.

Member 0 of the result gets the low-order bits of the multiplication.

Member 1 of the result gets the high-order bits of the multiplication.

Ī	5	151	<id>></id>	Result <id></id>	< <i>id</i> >	<id></id>
			Result Type		Operand 1	Operand 2

OpSMulExtended

Result is the full value of the signed integer multiplication of *Operand 1* and *Operand 2*.

Result Type must be from OpTypeStruct. The struct must have two members, and the two members must be the same type. The member type must be a scalar or vector of integer type.

Operand 1 and *Operand 2* must have the same type as the members of *Result Type*. These are consumed as signed integers.

Results are computed per component.

Member 0 of the result gets the low-order bits of the multiplication.

Member 1 of the result gets the high-order bits of the multiplication.

			1		
5	152	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

3.32.14 Bit Instructions

OpShiftRightLogical

Shift the bits in Base right by the number of bits specified in Shift. The most-significant bits will be zero filled.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

Shift is consumed as an unsigned integer. The result is undefined if *Shift* is greater than the bit width of the components of *Base*.

Results are computed per component.

5	194	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Base	Shift	

OpShiftRightArithmetic

Shift the bits in *Base* right by the number of bits specified in *Shift*. The most-significant bits will be filled with the sign bit from *Base*.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

Shift is treated as unsigned. The result is undefined if Shift is greater than the bit width of the components of Base.

Results are computed per component.

İ	5	195	< <i>id</i> >	Result <id></id>	<id>></id>	<id>></id>
			Result Type		Base	Shift

OpShiftLeftLogical

Shift the bits in Base left by the number of bits specified in Shift. The least-significant bits will be zero filled.

Result Type must be a scalar or vector of integer type.

The type of each *Base* and *Shift* must be a scalar or vector of integer type. *Base* and *Shift* must have the same number of components. The number of components and bit width of the type of *Base* must be the same as in *Result Type*.

Shift is treated as unsigned. The result is undefined if Shift is greater than the bit width of the components of Base.

The number of components and bit width of *Result Type* must match those *Base* type. All types must be integer types.

5	196	<id>></id>	Result <id></id>	<id></id>	<id>></id>
		Result Type		Base	Shift

OpBitwiseOr

Result is 1 if either *Operand 1* or *Operand 2* is 1. Result is 0 if both *Operand 1* and *Operand 2* are 0.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type. The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same number of components as *Result Type*. They must have the same component width as *Result Type*.

	5	197	<id></id>	Result <id></id>	<id></id>	<id></id>	
			Result Type		Operand 1	Operand 2	

OpBitwiseXor

Result is 1 if exactly one of *Operand 1* or *Operand 2* is 1. Result is 0 if *Operand 1* and *Operand 2* have the same value.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type. The type of Operand 1 and Operand 2 must be a scalar or vector of integer type. They must have the same number of components as Result Type. They must have the same component width as Result Type.

5	198	<id>></id>	Result <id></id>	<id>></id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpBitwiseAnd

Result is 1 if both Operand 1 and Operand 2 are 1. Result is 0 if either Operand 1 or Operand 2 are 0.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type. The type of Operand 1 and Operand 2 must be a scalar or vector of integer type. They must have the same number of components as Result Type. They must have the same component width as Result Type.

		* 1			
5	199	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpNot

Complement the bits of *Operand*.

Results are computed per component, and within each component, per bit.

Result Type must be a scalar or vector of integer type.

Operand's type must be a scalar or vector of integer type. It must have the same number of components as *Result Type*. The component width must equal the component width in *Result Type*.

Ī	4	200	< <i>id</i> >	Result <id></id>	<id></id>
			Result Type		Operand

OpBit	FieldIn	sert				Capability:			
Make a	а сору (of an object, with	a modified bit fiel	d that comes fron	n another object.	Shader			
Results	s are co	mputed per comp							
Result	Result Type must be a scalar or vector of integer type.								
The ty	pe of B	ase and Insert mu	st be the same as	Result Type.					
	Any result bits numbered outside [Offset, Offset + Count - 1] (inclusive) will come from the corresponding bits in Base.								
		s numbered in [<i>O</i> ₃] [0, <i>Count</i> - 1] of	ffset, Offset + Cou Insert.	<i>unt</i> - 1] come, in o	order, from the				
Insert.	It will		calar. <i>Count</i> is the n unsigned value.						
	Offset must be an integer type scalar. Offset is the lowest-order bit of the bit field. It will be consumed as an unsigned value.								
	The resulting value is undefined if <i>Count</i> or <i>Offset</i> or their sum is greater than the								
		s in the result.							
7	201	<id>></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id>></id>	<id>></id>		
		Result Type		Base	Insert	Offset	Count		

OpBitFieldS	Extract			Capability:	
Extract a bit	field from an object, wi	Shader			
Results are c	omputed per componen				
Result Type	nust be a scalar or vector	or of integer type.			
The type of	Base must be the same a	as Result Type.			
<i>Count</i> - 1] (i	eater than 0: The bits of clusive) become the big bits of the result will				
from Base. I	e an integer type scalar will be consumed as a te result will be 0.				
	e an integer type scalar et from <i>Base</i> . It will be	•••			
	value is undefined if <i>C</i> ber of bits in the result.				
6 202	<id>></id>	Result <id></id>	<id>></id>	<id></id>	< <i>id</i> >
	Result Type		Base	Offset	Count

OpBitl	OpBitFieldUExtract				Capability:	
The ser	mantics a	d from an object, with the same as with the control of the control	OpBitFieldSExtract v	with the exception	Shader	
6	203	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Base	Offset	Count

OpBitRev	erse			Capability: Shader
Reverse the	e bits in an ol	2.44.401		
Results are	computed po			
Result Type	e must be a se	calar or vector of integer t	ype.	
The type of	f <i>Base</i> must b			
	mber n of the th is the $\frac{OpT}{t}$	Base,		
4	204	<id></id>	Result <id></id>	< <i>id</i> >
		Result Type		Base

OpBitCount

Count the number of set bits in an object.

Results are computed per component.

Result Type must be a scalar or vector of integer type. The components must be wide enough to hold the unsigned *Width* of *Base* as an unsigned value. That is, no sign bit is needed or counted when checking for a wide enough result width.

Base must be a scalar or vector of integer type. It must have the same number of components as Result Type.

The result is the unsigned value that is the number of bits in *Base* that are 1.

4	205	< <i>id</i> >	Result <id></id>	<id></id>	
		Result Type		Base	

3.32.15 Relational and Logical Instructions

Result is **true** if any component of *Vector* is **true**, otherwise result is **false**. Result Type must be a Boolean type scalar. Vector must be a vector of Boolean type. 4 | 154 | <id> | Result <id> | <id> | <id> |

Vector

Result Type

OpAll						
Result is true if all components of <i>Vector</i> are true , otherwise result is false .						
Result T	<i>Type</i> must	be a Boolean type	scalar.			
Vector 1	nust be a	vector of Boolean	type.			
4	155	<id> Result Type</id>	Result <id></id>	<id></id>		
		Result Type		Vector		

OpIsNan

Result is **true** if *x* is an IEEE NaN, otherwise result is **false**.

Result Type must be a scalar or vector of Boolean type.

x must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

Results are computed per component.

Tresums and	results are compared per component					
4	156	< <i>id</i> >	Result <id></id>	< <i>id</i> >		
		Result Type		x		

OpIsInf

Result is **true** if x is an IEEE Inf, otherwise result is **false**

Result Type must be a scalar or vector of Boolean type.

x must be a scalar or vector of floating-point type. It must have the same number of components as Result Type.

4	157	<id>></id>	Result <id></id>	<id></id>
		Result Type		X

OpIsFinite	e	Capability:		
Result is tr	rue if x is an	calar or vector of Boo	therwise result is false . lean type. ype. It must have the same number of	Kernel
Results are	computed p	er component.		
4	158	<id></id>	Result <id></id>	<id>></id>
		Result Type		x

OpIsNorn	nal			Capability:
	rue if x is an a	Kernel		
	a scalar or vects as <i>Result T</i>			
Results are	computed po			
4	159	< <i>id</i> >	Result <id></id>	<id></id>
		Result Type		x

OpSignBitSet				Capability:
Result is true Result Type r x must be a s components	e if x has its must be a so scalar or vec as Result Ty	ppe.		Kernel
Results are computed per component.				
4	160	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		x

OpLess(OpLessOrGreater				
Result Ty x must be componed	true if x < wpe must b e a scalar o ents as Res	< y or $x > y$, where IEEE of e a scalar or vector of Booton vector of floating-point ult $Type$.	olean type.		Capability: Kernel
	y must have the same type as x. Results are computed per component.				
5	161	<id>></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		x	у

OpOrde	OpOrdered					
Result is result is Result Ty	Capability: Kernel					
x must be a scalar or vector of floating-point type. It must have the same number of components as <i>Result Type</i>.y must have the same type as x.						
Results are computed per component.						
5	162	<id>></id>	Result <id></id>	<id></id>	<id></id>	
		Result Type		x	у	

OpUnordered	Capability:			
Result is true if e	Kernel			
Result Type must				
x must be a scalar components as Re				
y must have the sa				
Results are compo				
5 163	<id>></id>	Result <id></id>	< <i>id</i> >	<id>></id>
	Result Type		X	у

OpLogicalEqual

Result is **true** if *Operand 1* and *Operand 2* have the same value. Result is **false** if *Operand 1* and *Operand 2* have different values.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

Results are computed per component.

5	164	<id>></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpLogicalNotEqual

Result is **true** if *Operand 1* and *Operand 2* have different values. Result is **false** if *Operand 1* and *Operand 2* have the same value.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

Results are computed per component.

5	165	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpLogicalOr

Result is **true** if either *Operand 1* or *Operand 2* is **true**. Result is **false** if both *Operand 1* and *Operand 2* are **false**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

	I	- I - I			
5	166	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpLogicalAnd

Result is **true** if both *Operand 1* and *Operand 2* are **true**. Result is **false** if either *Operand 1* or *Operand 2* are **false**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* must be the same as *Result Type*.

The type of *Operand 2* must be the same as *Result Type*.

Results are computed per component.

5	167	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpLogicalNot

Result is **true** if *Operand* is **false**. Result is **false** if *Operand* is **true**.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand* must be the same as *Result Type*.

Results are computed per component.

4	168	< <i>id</i> >	Result <id></id>	<id>></id>
		Result Type		Operand

OpSelect

Select between two objects.

Result Type must be a scalar or vector.

The type of *Object 1* must be the same as *Result Type*. *Object 1* is selected as the result if *Condition* is **true**.

The type of *Object 2* must be the same as *Result Type*. *Object 2* is selected as the result if *Condition* is **false**.

Condition must be a scalar or vector of Boolean type. It must have the same number of components as Result Type.

6	169	<id></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Condition	Object 1	Object 2

OpIEqual

Integer comparison for equality.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	170	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpINotEqual

Integer comparison for inequality.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

						i i
5	171	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	i
		Result Type		Operand 1	Operand 2	i

OpUGreaterThan

Unsigned-integer comparison if *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	172	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpSGreaterThan

Signed-integer comparison if *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

	5	173	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
			Result Type		Operand 1	Operand 2

OpUGreaterThanEqual

Unsigned-integer comparison if *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	174	<id>></id>	Result <id></id>	<id>></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpSGreaterThanEqual

Signed-integer comparison if *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	175	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpULessThan

Unsigned-integer comparison if *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	176	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpSLessThan

Signed-integer comparison if *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

	5	177	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
			Result Type		Operand 1	Operand 2

OpULessThanEqual

Unsigned-integer comparison if *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	178	<id></id>	Result <id></id>	<id></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpSLessThanEqual

Signed-integer comparison if *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of integer type. They must have the same component width, and they must have the same number of components as *Result Type*.

Results are computed per component.

		1 1				i
5	179	< <i>id</i> >	Result <id></id>	<id></id>	< <i>id</i> >	i
		Result Type		Operand 1	Operand 2	i

OpFOrdEqual

Floating-point comparison for being ordered and equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	180	<id>></id>	Result <id></id>	<id>></id>	< <i>id</i> >
		Result Type		Operand 1	Operand 2

OpFUnordEqual

Floating-point comparison for being unordered or equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

5	181	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpFOrdNotEqual

Floating-point comparison for being ordered and not equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	182	<id></id>	Result <id></id>	<id></id>	<id></id>	
		Result Type		Operand 1	Operand 2	

OpFUnordNotEqual

Floating-point comparison for being unordered or not equal.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	183	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpFOrdLessThan

Floating-point comparison if operands are ordered and *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	184	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpFUnordLessThan

Floating-point comparison if operands are unordered or *Operand 1* is less than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

roman rest to a roman rest rest rest rest rest rest rest rest						
	5	185	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >
			Result Type		Operand 1	Operand 2

OpFOrdGreaterThan

Floating-point comparison if operands are ordered and *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	186	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpFUnordGreaterThan

Floating-point comparison if operands are unordered or *Operand 1* is greater than *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

- 1							
	5	187	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
			Result Type		Operand 1	Operand 2	

${\bf OpFOrdLessThan Equal}$

Floating-point comparison if operands are ordered and *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	188	<id>></id>	Result <id></id>	<id>></id>	<id>></id>
		Result Type		Operand 1	Operand 2

OpFUnordLessThanEqual

Floating-point comparison if operands are unordered or *Operand 1* is less than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

5	189	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

OpFOrdGreaterThanEqual

Floating-point comparison if operands are ordered and *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

Results are computed per component.

5	190	<id></id>	Result <id></id>	<id></id>	<id></id>
		Result Type		Operand 1	Operand 2

OpFUnordGreaterThanEqual

Floating-point comparison if operands are unordered or *Operand 1* is greater than or equal to *Operand 2*.

Result Type must be a scalar or vector of Boolean type.

The type of *Operand 1* and *Operand 2* must be a scalar or vector of floating-point type. They must have the same type, and they must have the same number of components as *Result Type*.

5	191	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	
		Result Type		Operand 1	Operand 2	

3.32.16 Derivative Instructions

OpDPdx		Capability:		
	t as either <mark>O</mark> p kternal factor	Shader		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	f P must be tl			
This instru	ction is only			
4	207	<id></id>	Result <id></id>	<id></id>
		Result Type		P

OpDPdy		Capability:		
	t as either Op sternal factor	Shader		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	P must be tl			
This instruc	ction is only			
4	208	<id>></id>	Result <id></id>	<id>></id>
		Result Type		P

OpFwidth	OpFwidth					
Result is the san OpDPdy on <i>P</i> .	Shader					
Result Type mus	st be a sc	alar or vector of floating-point	type.			
The type of <i>P</i> m	The type of <i>P</i> must be the same as <i>Result Type</i> . <i>P</i> is the value to take the derivative of.					
This instruction						
4 209	9	<id></id>	Result <id></id>	<id></id>		
		Result Type		P		

OpDPdxFi	ine	Capability:		
	e partial deri encing based).	DerivativeControl		
Result Type	must be a so	calar or vector of floating-point	type.	
	P must be the			
This instruc	ction is only			
4	210	<id>></id>	Result <id></id>	<id></id>
		Result Type		P

OpDPdyFi	ine	Capability:		
	e partial deri encing based	DerivativeControl		
Result Type	must be a so	calar or vector of floating-point	type.	
The type of	P must be the			
This instruc	ction is only			
4	211	<id></id>	Result <id></id>	<id></id>
		Result Type		P

OpFwidth	Fine	Capability:		
Result is th		mputing the sum of the absolut	re values of OpDPdxFine and	DerivativeControl
Result Type	must be a se	calar or vector of floating-point	type.	
The type of	f P must be the			
This instruc	ction is only			
4	212	<id>></id>	Result <id></id>	<id>></id>
		Result Type		P

OpDPdxCoarse			Capability:		
	DerivativeControl				
Result is the partial de	rivative of P with respect to the v	vindow x coordinate. Will use			
local differencing base	ed on the value of P for the currer	nt fragment's neighbors, and			
will possibly, but not i	necessarily, include the value of <i>P</i>	for the current fragment.			
That is, over a given a	rea, the implementation can comp	oute x derivatives in fewer			
unique locations than	would be allowed for OpDPdxFir	ne.			
	_				
Result Type must be a	scalar or vector of floating-point	type.			
The type of <i>P</i> must be	the same as Result Type. P is the	value to take the derivative of.			
This instruction is only					
4 213	< <i>id></i>	Result <id></id>	<id></id>		
	Result Type		P		

OpDPdyCoarse			Capability:	
		DerivativeControl		
Result is the	partial deri	vative of <i>P</i> with respect to the w	vindow y coordinate. Will use	
local differen	cing based	on the value of P for the current	t fragment's neighbors, and	
will possibly,	, but not ne	cessarily, include the value of P	for the current fragment.	
That is, over	a given are	a, the implementation can comp	oute y derivatives in fewer	
unique location	ons than we	ould be allowed for OpDPdyFir	e.	
Result Type n	nust be a sc	ealar or vector of floating-point	type.	
The type of P	must be the			
This instructi	on is only			
4 2	214	< <i>id</i> >	Result <id></id>	< <i>id</i> >
		Result Type		P

OpFwidthC	Coarse	Capability:		
Result is the same as computing the sum of the absolute values of $OpDPdxCoarse$ and $OpDPdyCoarse$ on P .				DerivativeControl
Result Type r	must be a so			
The type of <i>I</i>	P must be the			
This instructi	ion is only			
4	215	<id></id>	Result <id></id>	<id></id>
		Result Type		P

3.32.17 Control-Flow Instructions

OpPhi

The SSA phi function.

The result is selected based on control flow: If control reached the current block from *Parent i*, *Result Id* gets the value that *Variable i* had at the end of *Parent i*.

Result Type can be any type.

Operands are a sequence of pairs: (*Variable 1, Parent 1* block), (*Variable 2, Parent 2* block), . . . Each *Parent i* block is the label of an immediate predecessor in the CFG of the current block. A *Parent i* block must not appear more than once in the operand sequence. All *Variables* must have a type matching *Result Type*.

Within a block, this instruction must appear before all non-**OpPhi** instructions (except for **OpLine**, which can be mixed with **OpPhi**).

3 + variable	245	< <i>id</i> >	Result <id></id>	< <i>id</i> >, < <i>id</i> >,
		Result Type		Variable, Parent,

OpLoopMerge

Declare a structured loop.

This instruction must immediately precede either an OpBranch or OpBranchConditional instruction. That is, it must be the second-to-last instruction in its block.

Merge Block is the label of the merge block for this structured loop.

Continue Target is the label of a block targeted for processing a loop "continue".

See Structured Control Flow for more detail.

İ	4	246	<id>></id>	<id>></id>	Loop Control
			Merge Block	Continue Target	

OpSelectionMerge

Declare a structured selection.

This instruction must immediately precede either an OpBranchConditional or OpSwitch instruction. That is, it must be the second-to-last instruction in its block.

Merge Block is the label of the merge block for this structured selection.

See Structured Control Flow for more detail.

3	247	<id></id>	Selection Control
		Merge Block	

OpLabel

The block label instruction: Any reference to a block is through the *Result* < *id*> of its label.

Must be the first instruction of any block, and appears only as the first instruction of a block.

2	248	Result <id></id>

OpBranch

Unconditional branch to Target Label.

Target Label must be the *Result <id>* of an OpLabel instruction in the current function.

This instruction must be the last instruction in a block.

2	249	<id>></id>
		Target Label

OpBranchConditional

If Condition is **true**, branch to True Label, otherwise branch to False Label.

Condition must be a Boolean type scalar.

True Label must be an OpLabel in the current function.

False Label must be an OpLabel in the current function.

Branch weights are unsigned 32-bit integer literals. There must be either no Branch Weights or exactly two branch weights. If present, the first is the weight for branching to True Label, and the second is the weight for branching to False Label. The implied probability that a branch is taken is its weight divided by the sum of the two Branch weights.

This instruction must be the last instruction in a block.

4 + variable	250	<id></id>	<id></id>	<id></id>	Literal, Literal,
		Condition	True Label	False Label	Branch weights

OpSwitch

Multi-way branch to one of the operand label $\langle id \rangle$.

Selector must have a type of OpTypeInt. Selector will be compared for equality to the Target literals.

Default must be the < id > of a label. If *Selector* does not equal any of the *Target* literals, control flow will branch to the *Default* label < id >.

Target must be alternating scalar integer *literals* and the *<id>>* of a label. If *Selector* equals a *literal*, control flow will branch to the following *label <id>>*. It is invalid for any two *literal* to be equal to each other. If *Selector* does not equal any *literal*, control flow will branch to the *Default* label *<id>>*. Each *literal* is interpreted with the type of *Selector*: The bit width of *Selector's* type will be the width of each *literal's* type. If this width is not a multiple of 32-bits, the literals must be sign extended when the OpTypeInt *Signedness* is set to 1. (See Literal Number.)

This instruction must be the last instruction in a block.

3 + variable	251	<id>> Selector</id>	<id> Default</id>	literal, label <id>, literal, label <id>,</id></id>
				 Target

OpKill	Capability:
	Shader
Fragment-shader discard.	
Ceases all further processing in any invocation that executes it: Only instructions these invocations executed before OpKill will have observable side effects. If this instruction is executed in non-uniform control flow, all subsequent control flow is non-uniform (for invocations that continue to execute).	
This instruction must be the last instruction in a block.	
This instruction is only valid in the Fragment Execution Model.	
1	252

OpReturn			
Return with no value from a fund	ction with void return type.		
This instruction must be the last instruction in a block.			
1	253		

OpReturnValue

Return a value from a function.

Value is the value returned, by copy, and must match the *Return Type* operand of the OpTypeFunction type of the OpFunction body this return instruction is in.

This instruction must be the last instruction in a block.

2	254	< <i>id</i> >
		Value

OpUnreachable

Pointer

Declares that this block is not reachable in the CFG.

This instruction must be the last instruction in a block.

1 255

Size

OpLifetimeSt	o p		Capability:
Declare that an	object is dead a	fter this instruction.	Kernel
<i>Pointer</i> is a point	nter to the object	et whose lifetime is ending. Its type must	
_	•	age Class Function.	
capability is no of memory wh	ot being used. If one of the being used if etime is er	ointer to a non-void type or the Addresses <i>Size</i> is non-zero, it is the number of bytes adding. Its type must be an integer type if its type has <i>Signedness</i> of 1, its sign bit	
3	257	< <i>id></i>	Literal Number
		Pointer	Size

3.32.18 Atomic Instructions

OpAtomicLoad

Atomically load through *Pointer* using the given *Semantics*. All subparts of the value that is loaded will be read atomically with respect to all other atomic accesses to it within *Scope*.

Result Type must be a scalar of integer type or floating-point type.

Pointer is the pointer to the memory to read. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

6	227	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory
		Result Type		Pointer	Scope	Semantics <id></id>
						Semantics

OpAtomicStore

Atomically store through *Pointer* using the given *Semantics*. All subparts of *Value* will be written atomically with respect to all other atomic accesses to it within *Scope*.

Pointer is the pointer to the memory to write. The type it points to must be a scalar of integer type or floating-point type.

Value is the value to write. The type of Value and the type pointed to by Pointer must be the same type.

5	228	< <i>id</i> >	Scope <id></id>	Memory Semantics	< <i>id</i> >	
		Pointer	Scope	<id>></id>	Value	
				Semantics		

OpAtomicExchange

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value from copying Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be a scalar of integer type or floating-point type.

7	229	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicCompareExchange

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by selecting Value if Original Value equals Comparator or selecting Original Value otherwise, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

Use Equal for the memory semantics of this instruction when Value and Original Value compare equal.

Use *Unequal* for the memory semantics of this instruction when *Value* and *Original Value* compare unequal. *Unequal* cannot be set to **Release** or **Acquire and Release**. In addition, *Unequal* cannot be set to a stronger memory-order then *Equal*.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*. This type must also match the type of *Comparator*.

		• •							
9	230	< <i>id</i> >	Result	< <i>id</i> >	Scope	Memory	Memory	< <i>id</i> >	< <i>id</i> >
		Result	<id></id>	Pointer	<id></id>	Semantics	Semantics	Value	Comparator
		Type			Scope	<id></id>	<id></id>		
					_	Equal	Unequal		

OpA	tomic	CompareExc	changeWeak					Capability: Kernel	
Atte	mpts to	o do the follow	wing:					Kerner	
				y with respec	t to any other	atomic access	es within		
		e same location							
		ough <i>Pointer</i> t							
		w Value by se							
_		alue otherwise							
3) st	ore the	New Value b	ack through <i>F</i>	Pointer.					
The	instruc	tion's result is							
The	weak c	compare-and-	when						
		-		•	can fail and s				
		idde equais ee igh <i>Pointer</i> .	imparator the	comparison	can ran and s	ore back the (Tiginai		
Value	c unou	ign i omer.							
Resu	ılt Type	must be an in	nteger type sc	alar.					
		0 1		0.1.1					
			ry semantics	of this instruc	ction when Val	ue and Origin	ial Value		
com	pare eq	luai.							
Use	Unequ	al for the mer	mory semanti	cs of this inst	ruction when	Value and Ori	ginal Value		
	-		•		se or Acquire		~		
		nnot be set to			_				
	•			•	•				
The	type of	f <i>Value</i> must b	e the same as	Result Type.	. The type of t	he value point	ed to by		
Poin	<i>ter</i> mu	st be the same	e as <i>Result Ty</i>	pe. This type	must also ma	tch the type of	f		
Com	parato	r.							
9	231	< <i>id</i> >	Result	< <i>id</i> >	Scope	Memory	Memory	< <i>id</i> >	< <i>id</i> >
		Result	<id>></id>	Pointer	<id></id>	Semantics	Semantics	Value	Comparator
		Туре			Scope	<id>></id>	<id></id>		
						Equal	Unequal		

OpAtomicIIncrement

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value through integer addition of 1 to Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

6	232	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory
		Result Type		Pointer	Scope	Semantics <id></id>
						Semantics

OpAtomicIDecrement

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value through integer subtraction of 1 from Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar. The type of the value pointed to by Pointer must be the same as Result Type.

6	233	<id></id>	Result <id></id>	<id></id>	Scope <id></id>	Memory
		Result Type		Pointer	Scope	Semantics <id></id>
						Semantics

OpAtomicIAdd

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by integer addition of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	234	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicISub

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by integer subtraction of Value from Original Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

7	235	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
					•	<id></id>	
						Semantics	

OpAtomicSMin

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by finding the smallest signed integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	236	< <i>id</i> >	Result <id></id>	<id></id>	Scope <id></id>	Memory	<id></id>
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicUMin

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by finding the smallest unsigned integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

7	237	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicSMax

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by finding the largest signed integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	238	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	<id></id>
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicUMax

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by finding the largest unsigned integer of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

t	7	239	∠; d>	Result <id></id>	< <i>id</i> >	Coope side	Mamari	<id>></id>
	/	239	<id></id>	Kesuit <iu></iu>	\ \ta>	Scope <id></id>	Memory	\(\iu\)
			Result Type		Pointer	Scope	Semantics	Value
			Tresuit Type		100000	Scope		,
							<1d>	
							Semantics	

OpAtomicAnd

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by the bitwise AND of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	240	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicOr

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through *Pointer* to get an *Original Value*,
- 2) get a New Value by the bitwise OR of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the Original Value.

Result Type must be an integer type scalar.

The type of *Value* must be the same as *Result Type*. The type of the value pointed to by *Pointer* must be the same as *Result Type*.

7	241	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicXor

Perform the following steps atomically with respect to any other atomic accesses within *Scope* to the same location:

- 1) load through Pointer to get an Original Value,
- 2) get a New Value by the bitwise exclusive OR of Original Value and Value, and
- 3) store the New Value back through Pointer.

The instruction's result is the *Original Value*.

Result Type must be an integer type scalar.

7	242	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory	< <i>id</i> >
		Result Type		Pointer	Scope	Semantics	Value
						<id></id>	
						Semantics	

OpAtomicFlag	TestAndSet			Capability:	
	the flag value points	Kernel			
	's result is true if the clear state immediate	· ·			
Result Type mu	st be a Boolean type				
Results are und	efined if an atomic f	lag is modified by	an instruction other		
than OpAtomic	FlagTestAndSet or (
6 318	< <i>id</i> >	Result <id></id>	< <i>id</i> >	Scope <id></id>	Memory
	Result Type		Pointer	Scope	Semantics <id></id>
					Semantics

OpAtomic	FlagClear	Capability:		
Pointer mu	sets the flag	Kernel		
	undefined if FlagTestAnd	Memory Semantics <id></id>		
7	317	<id> Pointer</id>	Scope <id>Scope</id>	Semantics

3.32.19 Primitive Instructions

OpEmitVertex	Capability:
	Geometry
Emits the current values of all output variables to the	
current output primitive. After execution, the values of	
all output variables are undefined.	
This instruction can only be used when only one stream	
is present.	
1	218

OpEndPrimitive	Capability:
	Geometry
Finish the current primitive and start a new one. No	
vertex is emitted.	
This instruction can only be used when only one	
stream is present.	
1	219

OpEmitStreamVertex	Capability:
	GeometryStreams
Emits the current values of all output variables	
to the current output primitive. After execution,	
the values of all output variables are undefined.	
Stream must be an <id> of a constant</id>	
instruction with a scalar integer type. That	
constant is the output-primitive stream number.	
This instruction can only be used when	
multiple streams are present.	
2 220	<id></id>
	Stream

Capability:
GeometryStreams
er.
<id>></id>
Stream

3.32.20 Barrier Instructions

OpControlBarrier

Wait for other invocations of this module to reach the current point of execution.

All invocations of this module within *Execution* scope must reach this point of execution before any invocation will proceed beyond it.

This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within *Execution*. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely.

If *Semantics* is non-zero (non-**None**), this instruction also serves as an **OpMemoryBarrier** instruction, and must also perform and adhere to the description and semantics of an **OpMemoryBarrier** instruction with the same *Memory* and *Semantics* operands. This allows atomically specifying both a control barrier and a memory barrier (that is, without needing two instructions). If *Semantics* is zero (**None**), *Memory* is ignored.

It is only valid to use this instruction with **TessellationControl**, **GLCompute**, or **Kernel** execution models.

4	224	Scope <id></id>	Scope <id></id>	Memory Semantics <id></id>
		Execution	Memory	Semantics

OpMemoryBarrier

Control the order that memory accesses are observed.

Ensures that memory accesses issued before this instruction will be observed before memory accesses issued after this instruction. This control is ensured only for memory accesses issued by this invocation and observed by another invocation executing within *Memory* scope.

Semantics declares what kind of memory is being controlled and what kind of control to apply.

To execute both a memory barrier and a control barrier, see OpControlBarrier.

3	225	Scope <id></id>	Memory Semantics <id></id>
		Memory	Semantics

3.32.21 Group Instructions

OpGroupAsyncCopy	7					Capability: Kernel	
Perform an asynchron Destination. The asyn	Kernei						
This instruction return the async copy to finis							
All invocations of this	module within	Execution mu	st reach this p	oint of execu	tion.		
This instruction is only control flow within <i>Ex</i> invocations will execu							
Result Type must be an	n OpTypeEvent	object.					
Destination must be a	pointer to a sca	lar or vector o	f floating-poir	nt type or inte	eger type.		
Destination pointer St	orage Class mus	st be Workgro	oup or Cross\	Vorkgroup.			
The type of Source mu	ist be the same	as <i>Destination</i>					
When <i>Destination</i> poi must be CrossWorkg from <i>Source</i> pointer.							
When <i>Destination</i> poil Class must be Workg each element to <i>Destin</i>	roup. In this cas						
Stride and NumElement Physical32 and 64 bit			•		•		
Event must be an OpT	ypeEvent.						
Event can be used to a shared by multiple cop	_			_	t to be		
If <i>Event</i> argument is n be returned.	ot OpConstantN	Null, the event	object supplie	ed in event arg	gument will		
9 259 <id> Result Type</id>	Result <id></id>	Scope <id> Execution</id>	<id> Destination</id>	<id> Source</id>	<id><id>NumElements</id></id>	<id> Stride</id>	<id> Event</id>

OpGroupWa	aitEvents			Capability:
		Kernel		
Wait for even	its generate	d by OpGroupAsync	Copy operations to complete. <i>Events</i>	
List points to performed.	Num Even			
All invocation	ns of this n	nodule within Executi	on must reach this point of execution.	
control flow v	within <i>Exe</i>	cution. This ensures the	orrectly if placed strictly within uniform hat if any invocation executes it, all re, an invocation may stall indefinitely.	
Execution mu	ıst be Wor	kgroup or Subgroup	Scope.	
Num Events r	must be a 3	2-bit integer type sca	lar.	
Events List m	nust be a po			
4 2	260	Scope <id></id>	< <i>id</i> >	<id></id>
		Execution	Num Events	Events List

OpGrou	OpGroupAll							
	Evaluates a predicate for all invocations in the group, resulting in true if predicate evaluates to true for all invocations in the group, otherwise the result is false .							
All invoc	cations of t	this module within Execu	tion must reach this poin	t of execution.				
flow with	nin <i>Execut</i>	only guaranteed to work of ion. This ensures that if a d elsewhere, an invocation	any invocation executes it					
Result Ty	pe must b	e a Boolean type.						
Executio	n must be	Workgroup or Subgrou	p Scope.					
Predicat	Predicate must be a Boolean type.							
5	5 261 < id> Result < id> Scope < id>							
		Result Type		Execution	Predicate			

OpGroupAny	Capability:							
Evaluates a predicate for all invocations in the to true for any invocation in the group, othe All invocations of this module within <i>Execut</i>	Groups							
This instruction is only guaranteed to work of flow within <i>Execution</i> . This ensures that if a execute it. If placed elsewhere, an invocation	any invocation executes it							
Result Type must be a Boolean type.								
Execution must be Workgroup or Subgrou	Execution must be Workgroup or Subgroup Scope.							
Predicate must be a Boolean type.								
5 262 < <i>id></i>	< <i>id</i> >							
Result Type		Execution	Predicate					

OpGr	oupBroa	dcast			Capability:	
1	n the <i>Value</i> utions in th	e of the invocation id	Groups			
All inv		of this module within	n <i>Execution</i> must r	each this point of		
uniform execute	m control	is only guaranteed to flow within <i>Execution</i> avocations will executive nitely.				
Result type so		st be a 32-bit or 64-b	it integer type or a	16, 32 or 64 float		
Execut	tion must	be Workgroup or S	ubgroup Scope.			
The ty	pe of Valu	ue must be the same	as Result Type.			
compo	onents or a	an integer datatype. A vector with 3 comp on the group.				
6	263	<id></id>	Result <id></id>	Scope <id></id>	< <i>id</i> >	<id></id>
		Result Type		Execution	Value	LocalId

OpGroupIAdd Capability: **Groups** An integer add group operation specified for all values of X specified by invocations in the group. The identity I is 0. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The type of *X* must be the same as *Result Type*. Result <id> Scope <id> Group Operation 264 <*id*> $\langle id \rangle$ Result Type Execution Operation X

OpGro	oupFAdd		Capability:			
		add group operation the group.	Groups			
The ide	entity I is	0.				
All inve		of this module within	n <i>Execution</i> must re	each this point of		
uniforn execute	n control	flow within Execution vocations will execu	on. This ensures that	placed strictly within at if any invocation where, an invocation		
Result	<i>Type</i> mus	at be a 16-bit, 32-bit,	or 64-bit floating-p	ooint type scalar.		
Executi	ion must	be Workgroup or S t				
The typ	e of X m	oust be the same as R				
6	265	<id></id>	Result <id></id>	Scope <id></id>	Group Operation	< <i>id></i>
		Result Type		Execution	Operation	X

OpGroupFMin Capability: **Groups** A floating-point minimum group operation specified for all values of Xspecified by invocations in the group. The identity I is +INF. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 16-bit, 32-bit, or 64-bit floating-point type scalar. Execution must be Workgroup or Subgroup Scope. The type of *X* must be the same as *Result Type*. Result <id> Scope <id> **Group Operation** 266 <*id*> $\langle id \rangle$ Result Type Execution Operation X

OpGro	oupUMir	1	Capability:			
	_	eger minimum group ocations in the group	Groups			
1	entity <i>I</i> is <i>I</i> is 64 bi	UINT_MAX when ts wide.				
All inve		of this module withi	n <i>Execution</i> must re	ach this point of		
uniforn execute	n control	flow within <i>Executi</i> nvocations will exec	on. This ensures that	placed strictly within t if any invocation where, an invocation		
Result	Type mus	at be a 32-bit or 64-b	it integer type scala	r.		
Executi	ion must	be Workgroup or S				
The typ	e of X m	oust be the same as R				
6	267	< <i>id</i> >	Result <id></id>	Scope <id></id>	Group Operation	<id>></id>
		Result Type		Execution	Operation	X

OpGroupSMin Capability: Groups A signed integer minimum group operation specified for all values of X specified by invocations in the group. The identity *I* is INT_MAX when *X* is 32 bits wide and LONG_MAX when *X* is 64 bits wide. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The type of *X* must be the same as *Result Type*. 268 <*id*> Result <id> Scope <id> **Group Operation** <*id*> Result Type Operation Execution X

OpGroupFMax			Capability:	
A floating-point maximum group oper specified by invocations in the group	Groups			
The identity <i>I</i> is -INF.				
All invocations of this module within execution.	Execution must re	ach this point of		
This instruction is only guaranteed to uniform control flow within <i>Execution</i> executes it, all invocations will execute may stall indefinitely.	n. This ensures that	t if any invocation		
Result Type must be a 16-bit, 32-bit,	or 64-bit floating-p	oint type scalar.		
Execution must be Workgroup or Su				
The type of X must be the same as R				
6 269 < <i>id</i> >	Result <id></id>	Scope <id></id>	Group Operation	<id>></id>
Result Type		Execution	Operation	X

OpGroupUMax Capability: Groups An unsigned integer maximum group operation specified for all values of Xspecified by invocations in the group. The identity I is 0. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be a 32-bit or 64-bit integer type scalar. Execution must be Workgroup or Subgroup Scope. The type of *X* must be the same as *Result Type*. 270 Result <id> Scope <id> **Group Operation** <*id*> $\langle id \rangle$ Result Type Execution Operation \mathbf{X}

OpGro	oupSMax	(Capability:			
_	_	r maximum group op ocations in the group	Groups			
1	entity <i>I</i> is its wide.	INT_MIN when <i>X</i> i				
All invo		of this module withi	n <i>Execution</i> must re	each this point of		
uniforn execute	n control	flow within Execution will executions will executions	on. This ensures the	placed strictly within at if any invocation where, an invocation		
X and F	Result Ty _l	pe must be a 32-bit of	or 64-bit OpTypeInt	data type.		
Executi	ion must	be Workgroup or S				
The typ	oe of X m	ust be the same as R				
6	271	<id> Result Type</id>	Result <id></id>	Scope <id> Execution</id>	Group Operation Operation	<id>X</id>

3.32.22 Device-Side Enqueue Instructions

OpEnqueueMark	Capability:									
Enqueue a marker command waits fo all previously enqu completes.	DeviceEnqueu	e								
	Result Type must be a 32-bit integer type scalar. A successful enqueue results in the value 0. A failed enqueue results in a non-0 value.									
Queue must be of	the type OpTy	peQueue.								
Num Events specif Events and must be integer.		•		•						
Wait Events specific OpTypeDeviceEve		vait event objects	s and must be a po	ointer to						
instruction. must b	Ret Event is a pointer to a device event which gets implicitly retained by this instruction. must be an OpTypePointer to OpTypeDeviceEvent. If Ret Event is set to null this instruction becomes a no-op.									
7 291 <ia< td=""><td>d></td><td>Result <id></id></td><td><<i>id</i>></td><td><<i>id</i>></td><td><id>></id></td><td><<i>id</i>></td></ia<>	d>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	<id>></id>	< <i>id</i> >				
Re	sult Type		Queue	Num Events	Wait Events	Ret Event				

OpEnqueueKernel

Capability:

Enqueue the function specified by *Invoke* and the NDRange specified by *ND Range*

for execution to the queue object specified by Queue.

Result Type must be a 32-bit integer type scalar. A successful enqueue results in the value 0. A failed enqueue results in a non-0 value.

Queue must be of the type OpTypeQueue.

Flags must be an integer type scalar. The content of Flags is interpreted as Kernel Enqueue Flags mask.

ND Range must be an OpTypeStruct created by OpBuildNDRange.

Num Events specifies the number of event objects in the wait list pointed by Wait *Events* and must be 32-bit integer type scalar, which is treated as unsigned integer.

Wait Events specifies the list of wait event objects and must be a pointer to OpTypeDeviceEvent.

Ret Event must be a pointer to OpTypeDeviceEvent which gets implicitly retained by this instruction.

Invoke must be an OpFunction whose OpTypeFunction operand has:

- Result Type must be OpTypeVoid.
- The first parameter must be an OpTypePointer to 8 bits OpTypeInt.
- Optional list of parameters that must be an OpTypePointer to the Workgroup Storage Class.

Param is the first parameter of the function specified by Invoke and must be a pointer to 8-bit integer type scalar.

Param Size is the size in bytes of the memory pointed by Param and must be a 32-bit integer type scalar, which is treated as unsigned integer.

Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as unsigned integer.

Each Local Size operand corresponds (in order) to one OpTypePointer to Workgroup Storage Class parameter to the Invoke function, and specifies the number of bytes of Workgroup storage used to back the pointer during the execution of the *Invoke* function.

1	3 +	292	! <id></id>	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >,
1	ari-		Result	<id></id>	Queue	Flags	ND	Num	Wait	Ret	Invoke	Param	Param	Param	<id>,</id>
a	ıble		Туре				Range	Events	Events	Event			Size	Align	
															Local
															Size

OpGetKernelNDrangeSubGroupCount Capability: **DeviceEnqueue** Returns the number of subgroups in each workgroup of the dispatch (except for the last in cases where the global size does not divide cleanly into work-groups) given the combination of the passed NDRange descriptor specified by ND Range and the function specified by Invoke. Result Type must be a 32-bit integer type scalar. ND Range must be an OpTypeStruct created by OpBuildNDRange. *Invoke* must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must be an OpTypePointer to 8 bits OpTypeInt. - Optional list of parameters that must be an OpTypePointer to the Workgroup Storage Class. Param is the first parameter of the function specified by Invoke and must be a pointer to 8-bit integer type scalar. Param Size is the size in bytes of the memory pointed by Param and must be a 32-bit integer type scalar, which is treated as unsigned integer. Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as unsigned integer. 293 <*id*> Result <id> <*id*> 8 <*id*> <*id*> <*id*> <*id*> Result Type ND Range Invoke Param Size Param Param Align

OpGetKer	nelNDrangeMa	xSubGroupSiz	e			Capability: DeviceEnque	eue			
	maximum sub-gpecified by ND									
Result Type	must be a 32-bi	t integer type so	calar.							
ND Range	must be an OpTy	peStruct create	d by OpBuildN	IDRange.						
	t be an OpFunction of the opFu		ypeFunction op	erand has:						
	arameter must b				kgroup Storage					
Class.										
	e first parameter r type scalar.	of the function	specified by In	woke and mus	t be a pointer to					
	Param Size is the size in bytes of the memory pointed by Param and must be a 32-bit integer type scalar, which is treated as unsigned integer.									
_	n is the alignmen									
8 294	unsigned integer	Result <id></id>	<id>></id>	< <i>id</i> >	< <i>id</i> >	<id>></id>	<id>></id>			
294	Result Type	Result \id>	ND Range	Invoke	Param	Param Size	Param			
	nesun 1ype		11D nunge	Invoice	1 arant	1 aranı bize	Align			

OpGetKerne	 IWorkGroupSize	е			Capability:				
	aximum work-growoke on the device	DeviceEnqueu	e						
Result Type m	ust be a 32-bit in	teger type scalar.							
- Result Type: - The first part - Optional list	Invoke must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must be an OpTypePointer to 8 bits OpTypeInt. - Optional list of parameters that must be an OpTypePointer to the Workgroup Storage Class.								
	irst parameter of integer type sca	•	ified by <i>Invoke</i> an	d must be a					
	Param Size is the size in bytes of the memory pointed by Param and must be a 32-bit integer type scalar, which is treated as unsigned integer.								
	s the alignment o								
7 295	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >			
	Result Type		Invoke	Param	Param Size	Param Align			

OpGetKernelPreferredWorkGroupSizeMultiple Capability: **DeviceEnqueue** Returns the preferred multiple of work-group size for the function specified by *Invoke*. This is a performance hint. Specifying a work-group size that is not a multiple of the value returned by this query as the value of the local work size will not fail to enqueue *Invoke* for execution unless the work-group size specified is larger than the device maximum. Result Type must be a 32-bit integer type scalar. *Invoke* must be an OpFunction whose OpTypeFunction operand has: - Result Type must be OpTypeVoid. - The first parameter must be an OpTypePointer to 8 bits OpTypeInt. - Optional list of parameters that must be an OpTypePointer to the Workgroup Storage Class. Param is the first parameter of the function specified by Invoke and must be a pointer to 8-bit integer type scalar. Param Size is the size in bytes of the memory pointed by Param and must be a 32-bit integer type scalar, which is treated as unsigned integer. Param Align is the alignment of Param and must be a 32-bit integer type scalar, which is treated as unsigned integer. <*id*> <id> <*id*> <*id*> 296 <*id*> Result <id> Result Type Invoke Param Param Size Param Align

OpRetainEvent	Capability:
Increments the reference count of the event object specified by <i>Event</i> .	DeviceEnqueue
Event must be an event that was produced by	
OpEnqueueKernel, OpEnqueueMarker or	
OpCreateUserEvent.	
2 297	<id></id>
	Event

OpReleaseEvent	Capability:
	DeviceEnqueue
Decrements the reference count of the event	
object specified by <i>Event</i> . The event object is	
deleted once the event reference count is zero,	
the specific command identified by this event	
has completed (or terminated) and there are no	
commands in any device command queue that	
require a wait for this event to complete.	
Event must be an event that was produced by	
OpEnqueueKernel, OpEnqueueMarker or	
OpCreateUserEvent.	
2 298	<id></id>
	Event

OpCreate	UserEvent	Capability: DeviceEnqueue	
event is se	t to a value of	e execution status of the created 2 (CL_SUBMITTED).	
Kesuu Typ	e must be Op	TypeDeviceEvent.	
3	299	Result <id></id>	
		Result Type	

OpIsVali	dEvent			Capability:
Returns to is false.	rue if the ev	DeviceEnqueue		
		Boolean type.		
Event mu	st be an Op'	ГуреDeviceEvent		
4	300	<id></id>	Result <id></id>	<id>></id>
		Result Type		Event

OpSetUserEv	entStatus	Capability:	
		DeviceEnqueue	
either 0 (CL_C	ion status of a us COMPLETE) to indexecution successor.		
OpCreateUserl	Event.	eEvent that was produced by	
Status must be	a 32-bit OpType		
3	301	< <i>id</i> >	< <i>id</i> >
		Event	Status

OpCaptureEventProfilingInfo Capability: DeviceEnqueue Captures the profiling information specified by Profiling Info for the command associated with the event specified by *Event* in the memory pointed by *Value*. The profiling information will be available in the memory pointed by Value once the command identified by Event has completed. Event must be an OpTypeDeviceEvent that was produced by OpEnqueueKernel or OpEnqueueMarker. Profiling Info must be an integer type scalar. The content of Profiling Info is interpreted as Kernel Profiling Info mask. Value must be a pointer to a scalar 8-bit integer type in the CrossWorkgroup Storage Class. When *Profiling Info* is **CmdExecTime**, *Value* must point to 128-bit memory range. The first 64 bits contain the elapsed time CL_PROFILING_COMMAND_END -CL PROFILING COMMAND START for the command identified by Event in nanoseconds. The second 64 bits contain the elapsed time CL_PROFILING_COMMAND_COMPLETE -CL_PROFILING_COMMAND_START for the command identified by Event in nanoseconds. Note: The behavior of this instruction is undefined when called multiple times for the same event. 302 <*id*> 4 <*id*> <id> Value Event Profiling Info

OpGetDefa	aultQueue		Capability:
		re queue. If a default device ed, a null queue object is	DeviceEnqueue
Result Type	must be an		
3	303	Result <id></id>	
		Result Type	

OpBuildNDRange

Given the global work size specified by GlobalWorkSize, local work size specified by LocalWorkSize and global work offset specified by GlobalWorkOffset, builds a 1D, 2D or 3D ND-range descriptor structure and returns it.

Result Type must be an OpTypeStruct with the following ordered list of members, starting from the first to last:

- 1) 32-bit integer type scalar, that specifies the number of dimensions used to specify the global work-items and work-items in the work-group.
- 2) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is **Physical64**. This member is an array of per-dimension unsigned values that describe the offset used to calculate the global ID of a work-item.
- 3) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is Physical64. This member is an array of per-dimension unsigned values that describe the number of global work-items in the dimensions that will execute the kernel function.
- 4) OpTypeArray with 3 elements, where each element is 32-bit integer type scalar when the addressing model is **Physical32** and 64-bit integer type scalar when the addressing model is **Physical64**. This member is an array of per-dimension unsigned values that describe the number of work-items that make up a work-group.

GlobalWorkSize must be a scalar or an array with 2 or 3 components. Where the type of each element in the array is 32-bit integer type scalar when the addressing model is **Physical32** or 64-bit integer type scalar when the addressing model is Physical64.

The type of *LocalWorkSize* must be the same as *GlobalWorkSize*.

The type of GlobalWorkOffset must be the same as GlobalWorkSize

Capability:

DeviceEnqueue

I ne typ	ic of Gibi	outworkojjset must	be the same as Gibbt	ii WOIKSIZE.		
6	304	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		GlobalWorkSize	LocalWorkSize	GlobalWorkOffset

3.32.23 Pipe Instructions

OpRe	adPipe					Capability:	
		from the pipe objuccessful and a ne	Pipes				
Result	t Type m	ust be a 32-bit int	eger type scalar.				
Pipe r	nust be a	an OpTypePipe w	ith ReadOnly ac	cess qualifier.			
	er must bge Class.		nter with the same	e data type as Pi	pe and a Generic		
		ust be a 32-bit into	eger type scalar t	hat represents th	e size in bytes of		
	0	ent must be a 32- ch packet in the pi	0 11	calar that presen	ts the alignment		
7	274	<id>></id>	Result <id></id>	<id>></id>	<id>></id>	<id>></id>	<id>></id>
		Result Type		Pipe	Pointer	Packet Size	Packet Alignment

OpW	ritePipe	;				Capability:	
		t from <i>Pointer</i> to accessful and a ne	Pipes				
Result	<i>t Type</i> m	ust be a 32-bit in	teger type scalar.				
Pipe r	must be	an OpTypePipe w	rith WriteOnly ac	ccess qualifier.			
	er must bge Class.		nter with the same	e data type as <i>Pip</i>	e and a Generic		
		ust be a 32-bit int the pipe	eger type scalar t	hat represents the	size in bytes of		
Packe	t Alignm	ent must be a 32-	-bit integer type s	calar that present	s the alignment		
	_	ch packet in the p					
7	275	<id>></id>	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Pipe	Pointer	Packet Size	Packet
							Alignment

OpRe	eserve	dReadPipe						Capability: Pipes	
Read specif from othery	fied by 0 <i>I</i>	Tipes							
Result	t Type	must be a 32	-bit integer t	ype scalar.					
Pipe 1	must b	e an OpType	Pipe with Re	eadOnly acc	ess qualifier.				
Reser	ve Id	must be an O	pTypeReserv	veId.					
Index	must	be a 32-bit in	iteger type so	alar, which	is treated as uns	signed value			
Pointe Class.		st be an OpTy	pePointer w	ith the same	data type as Pi	pe and a Ge	neric Storage		
		must be a 32 e pipe	-bit integer t	ype scalar th	at represents th	e size in byt	es of each		
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe								
9	276	<id>></id>	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result	<id></id>	Pipe	Reserve	Index	Pointer	Packet	Packet
		Туре			Id			Size	Alignment

OpReserv	edWritePipe						Capability Pipes	:
Write a pact pipe object from 0 otherwise.	Tipes							
Result Type	e must be a 32	2-bit integer	type scalar.					
Pipe must	be an OpType	Pipe with W	riteOnly ac	cess qualifier.				
Reserve Id	must be an O	pTypeReser	veId.					
Index must	be a 32-bit ir	nteger type so	calar, which i	is treated as un	signed value			
Pointer mu	st be an OpTy	ypePointer w	ith the same	data type as Pi	pe and a Ge	neric Storage		
Packet Size packet in the	e must be a 32 ne pipe	es of each						
_	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes							
	of each packet in the pipe							
9 277	< <i>id</i> >	Result	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >	<id>></id>	< <i>id</i> >
	Result	<id></id>	Pipe	Reserve	Index	Pointer	Packet	Packet
	Туре			Id			Size	Alignment

OpRe	eserveRo	eadPipePackets	Capability:				
		Packets entries fo id reservation ID	Pipes				
Result	<i>t Type</i> m	ust be an OpType	ReserveId.				
Pipe r	must be	an OpTypePipe w	ith ReadOnly acc	cess qualifier.			
Num I value.		must be a 32-bit ir	nteger type scalar,	which is treated	as unsigned		
		ust be a 32-bit into	eger type scalar th	nat represents the	size in bytes of		
Packe	t Alignm	ent must be a 32-	bit integer type so	calar that presents	the alignment		
	_	ch packet in the pi					
7	278	<id>></id>	< <i>id</i> >	<id></id>			
		Result Type	Num Packets	Packet Size	Packet Alignment		
							1111gmmem

OpRo	OpReserveWritePipePackets						
	Reserve <i>num_packets</i> entries for writing to the pipe object specified by <i>Pipe</i> . Result is a valid reservation ID if the reservation is successful.						
Pipe 1	must be	an OpTypePipe v	with WriteOnly a	ccess qualifie	r.		
Num	Packets 1	must be a 32-bit	OpTypeInt which	is treated as u	insigned value.		
Resul	<i>t Type</i> m	ust be an OpTyp	eReserveId.				
		ust be a 32-bit in the pipe	teger type scalar	that represents	s the size in bytes of		
	_	nent must be a 32 th packet in the p		scalar that pre	sents the alignment		
7	279	< <i>id></i>	< <i>id</i> >	< <i>id</i> >			
		Result Type	Packet Size	Packet			
							Alignment

OpCom	mitReadF	Pipe			Capability:		
	Indicates that all reads to <i>Num Packets</i> associated with the reservation specified by <i>Reserve Id</i> and the pipe object specified by <i>Pipe</i> are completed.						
Pipe mus	st be an O	pTypePipe with ReadOn	ly access qualifier.				
Reserve .	<i>Id</i> must be	e an OpTypeReserveId.					
	<i>ize</i> must be the pipe	e a 32-bit integer type sca	alar that represents the size	ze in bytes of each			
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe						
5	5 280 < <i>id></i> < <i>id></i>						
		Pipe	Reserve Id	Packet Size	Packet Alignment		

OpCom	mitWrite	Pipe			Capability:		
	Indicates that all writes to <i>Num Packets</i> associated with the reservation specified by <i>Reserve Id</i> and the pipe object specified by <i>Pipe</i> are completed.						
Pipe mu	st be an O	pTypePipe with W	riteOnly access qualifier.				
Reserve	Id must be	e an OpTypeReserv	veId.				
	<i>ize</i> must be the pipe	e a 32-bit integer t	ype scalar that represents th	ne size in bytes of each			
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe						
5	5 281 <id> <id> <id> </id></id></id>						
		Pipe	Reserve Id	Packet Size	Packet Alignment		

Return tr Result Ty	idReservel rue if Reser pe must be	Capability: Pipes		
Reserve I	d must be	an OpTypeReserveId.		
4	282	<id></id>	Result <id></id>	<id>></id>
		Result Type		Reserve Id

OpGet	NumPip	ePackets			Capability:	
			Pipes			
Result	is the nur	nber of available ent	ries in the pipe obj	ect specified by Pipe.		
		available entries in a be considered imme		value. The value		
1	<i>Type</i> mused value.	st be a 32-bit integer	type scalar, which	should be treated as		
Pipe m qualifie		OpTypePipe with R	eadOnly or Write	eOnly access		
Packet	Size mus	t be a 32-bit integer	type scalar that ren	resents the size in		
		cket in the pipe	71			
	•					
Packet	Alignmer	nt must be a 32-bit in	nteger type scalar t	hat presents the		
alignm	ent in by	tes of each packet in				
6	283	< <i>id</i> >	Result <id></id>	<id>></id>	<id></id>	< <i>id</i> >
		Result Type		Pipe	Packet Size	Packet Alignment

OpGet	MaxPip	ePackets	Capability:			
		ximum number of paragraphs e was created.	Pipes			
	<i>Type</i> mused value.	at be a 32-bit integer	type scalar, which sh	nould be treated as		
Pipe m qualifie		OpTypePipe with R	eadOnly or WriteO	only access		
1		t be a 32-bit integer to cket in the pipe	type scalar that repre	esents the size in		
Packet	Alignmer	at must be a 32-bit in	teger type scalar tha	t presents the		
1	-	tes of each packet in		ı		
6	284	< <i>id</i> >	Result <id></id>	< <i>id</i> >	< <i>id</i> >	< <i>id</i> >
		Result Type		Pipe	Packet Size	Packet Alignment

OpGroupReserveReadPipePackets Capability: **Pipes** Reserve Num Packets entries for reading from the pipe object specified by Pipe at group level. Result is a valid reservation id if the reservation is successful. The reserved pipe entries are referred to by indices that go from 0... Num Packets - 1. All invocations of this module within *Execution* must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Result Type must be an OpTypeReserveId. Execution must be Workgroup or Subgroup Scope. Pipe must be an OpTypePipe with ReadOnly access qualifier. Num Packets must be a 32-bit integer type scalar, which is treated as unsigned value. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe 285 <*id*> Result <id> Scope <id> <*id*> $\langle id \rangle$ $\langle id \rangle$ <*id*> Result Type Execution Pipe Num Packet Size Packet **Packets** Alignment

OpG	roupR	eserveWritePi _j	pePackets				Capability:	
Reserve <i>Num Packets</i> entries for writing to the pipe object specified by <i>Pipe</i> at group level. Result is a valid reservation ID if the reservation is successful.								
The r	eserved	l pipe entries ar	e referred to by	indices that go	from 0 <i>Ni</i>	ım Packets - 1.		
All ir	vocatio	ons of this mode	ule within Exec	ution must reac	h this point of	execution.		
contr	ol flow	within Execution	ranteed to work on. This ensured If placed elsewl	s that if any inv	ocation execut	es it, all		
Resul	lt Type	must be an OpT	TypeReserveId.					
Execu	ution m	ust be Workgr	oup or Subgro	up Scope.				
Pipe	must be	e an OpTypePip	e with WriteO	nly access qual	ifier.			
Num	Packets	must be a 32-t	oit integer type	scalar, which is	treated as uns	igned value.		
	Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe							
	Packet Alignment must be a 32-bit integer type scalar that presents the alignment in							
_	bytes of each packet in the pipe							
8	286	< <i>id</i> >	Result <id></id>	Scope <id></id>	< <i>id</i> >	< <i>id</i> >	<id></id>	<id>></id>
		Result Type		Execution	Pipe	Num	Packet Size	Packet
						Packets		Alignment

OpGroupCommitReadPipe Capability: **Pipes** A group level indication that all reads to Num Packets associated with the reservation specified by Reserve Id to the pipe object specified by Pipe are completed. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Execution must be Workgroup or Subgroup Scope. Pipe must be an OpTypePipe with ReadOnly access qualifier. Reserve Id must be an OpTypeReserveId. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe 287 Scope <id> 6 <*id*> <*id*> <*id*> <*id*> Execution Pipe Reserve Id Packet Size Packet Alignment

OpGroupCommitWritePipe Capability: **Pipes** A group level indication that all writes to Num Packets associated with the reservation specified by Reserve Id to the pipe object specified by Pipe are completed. All invocations of this module within Execution must reach this point of execution. This instruction is only guaranteed to work correctly if placed strictly within uniform control flow within Execution. This ensures that if any invocation executes it, all invocations will execute it. If placed elsewhere, an invocation may stall indefinitely. Execution must be Workgroup or Subgroup Scope. Pipe must be an OpTypePipe with WriteOnly access qualifier. Reserve Id must be an OpTypeReserveId. Packet Size must be a 32-bit integer type scalar that represents the size in bytes of each packet in the pipe Packet Alignment must be a 32-bit integer type scalar that presents the alignment in bytes of each packet in the pipe 288 Scope <id> 6 <*id*> <*id*> <*id*> <*id*> Execution Pipe Reserve Id Packet Size Packet Alignment

A Changes

A.1 Changes from Version 0.99, Revision 31

- Added the PushConstant Storage Class.
- Added OpIAddCarry, OpISubBorrow, OpUMulExtended, and OpSMulExtended.
- Added OpInBoundsPtrAccessChain.
- Added the Decoration NoContraction to prevent combining multiple operations into a single operation (bug 14396).
- Added sparse texturing (14486):
 - Added **OpImageSparse...** for accessing images that might not be resident.
 - Added **MinLod** functionality for accessing images with a minimum level of detail.
- Added back the **Alignment** Decoration, for the **Kernel** capability (14505).
- Added a NonTemporal Memory Access (14566).
- Structured control flow changes:
 - Changed structured loops to have a structured continue Target in OpLoopMerge (14422).
 - Added rules for how "fall through" works with **OpSwitch** (13579).
 - Added definitions for what is "inside" a structured control-flow construct (14422).
- Added **SubpassData** Dim to support input targets written by a previous subpass as an output target (14304). This is also a Decoration and a Capability, and can be used by some image ops to read the input target.
- Added OpTypeForwardPointer to establish the Storage Class of a forward reference to a pointer type (13822).
- · Improved Debuggability
 - Changed OpLine to not have a target <id>, but instead be placed immediately preceding the instruction(s) it is annotating (13905).
 - Added OpNoLine to terminate the affect of **OpLine** (13905).
 - Changed OpSource to include the source code:
 - * Allow multiple occurrences.
 - * Be mixed in with the OpString instructions.
 - * Optionally consume an OpString result to say which file it is annotating.
 - * Optionally include the source text corresponding to that OpString.
 - * Included adding OpSourceContinued for source text that is too long for a single instruction.
- Added a large number of Capabilities for subsetting functionality (14520, 14453), including 8-bit integer support for OpenCL kernels.
- Added VertexIndex and InstanceIndex BuiltIn Decorations (14255).
- Added GenericPointer capability that allows the ability to use the Generic Storage Class (14287).
- Added IndependentForwardProgress Execution Mode (14271).
- Added OpAtomicFlagClear and OpAtomicFlagTestAndSet instructions (14315).
- Changed OpentryPoint to take a list of **Input** and **Output** < id> for declaring the entry point's interface.
- · Fixed internal bugs
 - 14411 Added missing documentation for mad_sat OpenCL extended instructions (enums existed, just the documentation was missing)
 - 14241 Removed shader capability requirement from OpImageQueryLevels and OpImageQuerySamples.
 - 14241 Removed unneeded OpImageQueryDim instruction.

- 14241 Filled in TBD section for OpAtomicCompareExchangeWeek
- 14366 All OpSampledImage must appear before uses of sampled images (and still in the first block of the entry point).
- 14450 DeviceEnqueue capability is required for OpTypeQueue and OpTypeDeviceEvent
- 14363 OpTypePipe is opaque moved packet size and alignment to opcodes
- 14367 Float16Buffer capability clarified
- 14241 Clarified how OpSampledImage can be used
- 14402 Clarified OpTypeImage encodings for OpenCL extended instructions
- 14569 Removed mention of non-existent OpFunctionDecl
- 14372 Clarified usage of OpGenericPtrMemSemantics
- 13801 Clarified the **SpecId** Decoration is just for constants
- 14447 Changed literal values of Memory Semantic enums to match OpenCL/C++11 atomics, and made the Memory Semantic None and Relaxed be aliases
- 14637 Removed subgroup scope from OpGroupAsyncCopy and OpGroupWaitEvents

A.2 Changes from Version 0.99, Revision 32

- Added UnormInt101010_2 to the Image Channel Data Type table.
- Added place holder for C++11 atomic Consume Memory Semantics along with an explicit AcquireRelease memory semantic.
- Fixed internal bugs:
 - 14690 OpSwitch literal width (and hence number of operands) is determined by the type of Selector, and be rigorous about how sub-32-bit literals are stored.
 - 14485 The client API owns the semantics of built-ins that only have "pass through" semantics WRT SPIR-V.
- Fixed public bugs:
 - 1387 Don't describe result type of OpImageWrite.

A.3 Changes from Version 1.00, Revision 1

- Adjusted Capabilities:
 - Split geometry-stream functionality into its own **GeometryStreams** capability (14873).
 - Have InputAttachmentIndex to depend on InputAttachment instead of Shader (14797).
 - Merge AdvancedFormats and StorageImageExtendedFormats into just StorageImageExtendedFormats (14824).
 - Require StorageImageReadWithoutFormat and StorageImageWriteWithoutFormat to read and write storage images with an Unknown Image Format.
 - Removed the **ImageSRGBWrite** capability.
- · Clarifications
 - RelaxedPrecision Decoration can be applied to OpFunction (14662).
- Fixed internal bugs:
 - 14797 The literal argument was missing for the **InputAttachmentIndex** Decoration.
 - 14547 Remove the **FragColor** BuiltIn, so that no implicit broadcast is implied.
 - 13292 Make statements about "Volatile" be more consistent with the memory model specification (non-functional change).

- 14948 Remove image-"Query" overloading on image/sampled-image type and "fetch" on non-sampled images, by adding the OpImage instruction to get the image from a sampled image.
- 14949 Make consistent placement between **OpSource** and **OpSourceExtension** in the logical layout of a module.
- 14865 Merge WorkgroupLinearId with LocalInvocationId BuiltIn Decorations.
- 14806 Include 3D images for OpImageQuerySize.
- 14325 Removed the **Smooth Decoration**.
- 12771 Make the version word formatted as: "0 | Major Number | Minor Number | 0" in the physical layout.
- 15035 Allow OpTypeImage to use a *Depth* operand of 2 for not indicating a depth or non-depth image.
- 15009 Split the OpenCL Source Language into two: OpenCL_C and OpenCL_CPP.
- 14683 OpSampledImage instructions can only be the consuming block, for scalars, and directly consumed by an image lookup or query instruction.
- 14325 mutual exclusion validation rules of Execution Modes and Decorations
- 15112 add definitions for invocation, dynamically uniform, and uniform control flow.

· Renames

- InputTargetIndex Decoration → InputAttachmentIndex
- InputTarget Capability → InputAttachment
- InputTarget $Dim \rightarrow SubpassData$
- WorkgroupLocal Storage Class → Workgroup
- WorkgroupGlobal Storage Class \rightarrow CrossWorkgroup
- PrivateGlobal Storage Class \rightarrow Private
- OpAsyncGroupCopy → OpGroupAsyncCopy
- OpWaitGroupEvents → OpGroupWaitEvents
- InputTriangles Execution Mode → Triangles
- InputQuads Execution Mode → Quads
- InputIsolines Execution Mode → Isolines

A.4 Changes from Version 1.00, Revision 2

- Updated example at the end of Section 1 to conform to the KHR_vulkan_glsl extension and treat OpTypeBool as an abstract type.
- Adjusted Capabilities:
 - MatrixStride depends on Matrix (15234).
 - Sample, SampleId, SamplePosition, and SampleMask depend on SampleRateShading (15234).
 - ClipDistance and CullDistance BuiltIns depend on, respectively, ClipDistance and CullDistance (1407, 15234).
 - ViewportIndex depends on MultiViewport (15234).
 - AtomicCounterMemory should be the AtomicStorage (15234).
 - Float16 has no dependencies (15234).
 - Offset Decoration should only be for Shader (15268).
 - Generic Storage Class is supposed to need the GenericPointer Capability (14287).
 - Remove capability restriction on the **BuiltIn** Decoration (15248).
- Fixed internal bugs:
 - 15203 Updated description of SampleMask BuiltIn to include "Input or output...", not just "Input..."
 - 15225 Include no re-association as a constraint required by the NoContraction Decoration.
 - 15210 Clarify OpPhi semantics that operand values only come from parent blocks.

- 15239 Add OpImageSparseRead, which was missing (supposed to be 12 sparse-image instructions, but only 11 got incorporated, this adds the 12th).
- 15299 Move OpUndef back to the Miscellaneous section.
- 15321 OpTypeImage does not have a *Depth* restriction when used with **SubpassData**.
- 14948 Fix the **Lod** Image Operands to allow both integer and floating-point values.
- 15275 Clarify specific storage classes allowed for atomic operations under universal validation rules "Atomic access rules".
- 15501 Restrict **Patch** Decoration to one of the tessellation execution models.
- 15472 Reserved use of OpImageSparseSampleProjImplicitLod, OpImageSparseSampleProjExplicitLod, OpImageSparseSampleProjDrefImplicitLod, and OpImageSparseSampleProjDrefExplicitLod.
- 15459 Clarify what makes different aggregate types in "Types and Variables".
- 15426 Don't require OpQuantizeToF16 to preserve NaN patterns.
- 15418 Don't set both **Acquire** and **Release** bits in Memory Semantics.
- 15404 OpFunction Result <id> can only be used by OpFunctionCall, OpEntryPoint, and decoration instructions.
- 15437 Restrict element type for OpTypeRuntimeArray by adding a definition of concrete types.
- 15403 Clarify OpTypeFunction can only be consumed by OpFunction and functions can only return concrete and abstract types.
- Improved accuracy of the opcode word count in each instruction regarding which operands are optional. For sampling operations with explicit LOD, this included not marking the required LOD operands as optional.
- Clarified that when NonWritable, NonReadable, Volatile, and Coherent Decorations are applied to the Uniform storage class, the BufferBlock decoration must be present.
- · Fixed external bugs:
 - 1413 (see internal 15275)
 - 1417 Added definitions for block, dominate, post dominate, CFG, and back edge. Removed use of "dominator tree".

A.5 Changes from Version 1.00, Revision 3

• Added definition of derivative group, and use it to say when derivatives are well defined.