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```

1 Basic

1.1 Default Code [37e06b]

```
#include <bits/stdc++.h>
using namespace std;

#define iter(v) v.begin(),v.end()
#define SZ(v) int(v.size())
#define pb emplace_back
#define ff first
#define ss second

using ll = long long;
using pii = pair<int, int>;
using pll = pair<ll, ll>;

#ifdef zisk
void debug(){cerr << "\n";}
template<class T, class ... U>
```

```
void debug(T a, U ... b){cerr << a << " ", debug(b...)</pre>
template < class T> void pary(T 1, T r){
  while (1 != r) cerr << *1 << "
  cerr << "\n";</pre>
#else
#define debug(...) void()
#define pary(...) void()
#endif
template < class A, class B>
ostream& operator<<(ostream& o, pair<A,B> p)
{ return o << '(' << p.ff << ','
                                 << p.ss << ')'; }
int main(){
  ios_base::sync_with_stdio(0); cin.tie(0);
1.2 .vimrc [476d29]
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
map <F9> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -
    Wextra -Wshadow -O2 -Dzisk -g -fsanitize=undefined,
    address<CR>
map <F8> :!./%:r<CR>
inoremap {<CR> {<CR>}<ESC>ko
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
inoremap fj <ESC>
vnoremap fj <ESC>
" -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG
    Fast IO [c7ddfd]
// from JAW
inline int my_getchar() {
  const int N = 1 << 20;
  static char buf[N];
  static char *p = buf , *end = buf;
  if(p == end) {
    if((end = buf + fread(buf , 1 , N , stdin)) == buf)
         return EOF;
    p = buf;
  return *p++;
inline int readint(int &x) {
  static char c , neg;
  while((c = my_getchar()) < '-') {</pre>
    if(c == EOF) return 0;
  neg = (c == '-') ? -1 : 1;
  x = (neg == 1) ? c - '0' : 0;
  while((c = my_getchar()) >= '0') x = (x << 3) + (x <<
       1) + (c - '0');
  x *= neg;
  return 1;
const int kBufSize = 524288;
char inbuf[kBufSize];
char buf_[kBufSize]; size_t size_;
inline void Flush_() { write(1, buf_, size_); size_ =
    0; }
inline void CheckFlush_(size_t sz) { if (sz + size_ >
    kBufSize) Flush_(); }
inline void PutInt(int a) {
  static char tmp[22] = "01234567890123456789 \ n";
  CheckFlush_(10);
  if(a < 0){
    *(buf_ + size_) = '-';
    a = ~a + 1;
    size_++;
  int tail = 20;
  if (!a) {
    tmp[--tail] = '0';
  } else {
    for (; a; a /= 10) tmp[--tail] = (a % 10) ^ '0';
```

```
memcpy(buf_ + size_, tmp + tail, 21 - tail);
  size_ += 21 - tail;
int main(){
 Flush_();
  return 0;
1.4 Random [4cf9ed]
mt19937 rng(chrono::system_clock::now().
    time_since_epoch().count());
     PBDS Tree [7c702a]
```

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
using Tree = tree<int, null_type, less<>, rb_tree_tag,
    tree_order_node_statistics_update>;
// .find_by_order(x)
// .order_of_key(x)
```

1.6 Pragma [6006f6]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

1.7 SVG Writer [7adcc8]

```
class SVG {
   void p(string_view s) { o << s; }</pre>
   void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
      o << s.substr(0, i) << v, p(s.substr(i + 1), vs...)</pre>
   ofstream o; string c = "red";
public: // SVG svg("test.svg", 0, 0, 100, 100)
SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f
         ) {
      p("<svg xmlns='http://www.w3.org/2000/svg' "
         "viewBox='$ $ $ '>\n"
         "<style>*{stroke-width:0.5%;}</style>\n",
   x1, -y2, x2 - x1, y2 - y1); }
~SVG() { p("</svg>\n"); }
   void color(string nc) { c = nc; }
  void line(auto x1, auto y1, auto x2, auto y2) {
  p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n
        x1, -y1, x2, -y2, c); }
  void circle(auto x, auto y, auto r) {
  p("<circle cx='$' cy='$' r='$' stroke='$' "
    "fill='none'/>\n", x, -y, r, c); }

   void text(auto x, auto y, string s, int w = 12) {
 p("<text x='$' y='$' font-size='$px'>$</text>\n",
        x, -y, w, s); }
};
```

2 Data Structure

Heavy-Light Decomposition [f2dbca]

```
struct HLD{ // 1-based
  int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
  explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1), ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
```

```
}
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  }
  void build(){ dfs(1, 1), cut(1, 1); }
  int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // lca
};
```

2.2 Link Cut Tree [cf4f34]

```
// 1-based
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node() : pa{0}, ch{0}, 0}, rev{false}, size{1}, v{},
      sum{}, rsum{}, sv{}, sub{}, vir{} {}
  };
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is_rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false:
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
cur.rev ^= 1;
  void rotate(int u) {
    int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  void splay(int u) {
    vector<int> stk = {u};
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    up(u);
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
```

```
cur.vir = cur.vir + o[rc].sub - o[last].sub;
                                                             int Rank(node *o, 11 key) { // num of key < key</pre>
                                                                if (!o) return 0;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0;
    for (access(u); u; u = lc) down(la = u);
                                                               if (!o) return 0;
    return la;
                                                                  node *t = o;
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
                                                                  return 1;
                                                                }
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); 
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
                                                                node *a, *b;
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
    return find_root(u) == find_root(v); }
  void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
                                                             namespace kdt {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
                                                                point p[maxn];
#undef rc
};
2.3
     Treap [2ac37e]
mt19937 rng(880301);
struct node {
  11 data; int sz;
  node *1, *r;
  node(11 k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
                                                                  if (~lc[m]) {
  }
  void down() {}
node pool[1000010]; int pool_cnt = 0;
node *newnode(ll k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
                                                                  if (~rc[m]) {
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
                                                                  }
           a;
  return b->down(), b->l = merge(a, b->l), b->up(), b;
                                                                  return m;
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)</pre>
                                                                    return false;
   a = o, split(o \rightarrow r, a \rightarrow r, b, k), <math>a \rightarrow up();
                                                                  return true;
  else b = o, split(o->1, a, b->1, k), b->up();
// a: size k, b: size n - k
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) \leftarrow k) return a = o, b = 0, void();
                                                                void dfs(
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
node *kth(node *o, ll k) { // 1-based
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow l) - 1);
                                                                  } else {
```

```
if (o->data < key)</pre>
    return sz(o\rightarrow 1) + 1 + Rank(o\rightarrow r, key);
  else return Rank(o->1, key);
bool erase(node *&o, 11 k) {
  if (o->data == k) {
    o->down(), o = merge(o->1, o->r);
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, ll k) {
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
tuple<node*, node*, node*> interval(node *&o, int 1,
  int r) { // 1-based
node *a, *b, *c; // b: [l, r]
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
  return make_tuple(a, b, c);
2.4 KD Tree [375ca2]
  int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
    function<bool(const point &, const point &)> f =
      [dep](const point &a, const point &b) {
        if (dep & 1) return a.x < b.x;</pre>
        else return a.y < b.y;</pre>
    int m = (1 + r) >> 1;
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
    yl[m] = yr[m] = p[m].y;
    lc[m] = build(1, m, dep + 1);
      xl[m] = min(xl[m], xl[lc[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
  bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
        q.y < y1[o] - ds || q.y > yr[o] + ds
  long long dist(const point &a, const point &b) {
    return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
      const point &q, long long &d, int o, int dep = 0)
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x ||</pre>
        !(dep & 1) && q.y < p[o].y) {
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
```

```
National Taiwan University
       if (~rc[o]) dfs(q, d, rc[o], dep + 1);
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    }
  void init(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
    root = build(0, v.size());
  long long nearest(const point &q) {
    long long res = 1e18;
    dfs(q, res, root);
    return res;
} // namespace kdt
2.5 Leftist Tree [e91538]
struct node {
 11 v, data, sz, sum;
node *1, *r;
  node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->1)) swap(a->r, a->1);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a;
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
2.6 Convex 1D/1D [a449dd]
template < class T>
struct DynamicHull {
  struct seg { int x, l, r; };
  T f; int C; deque<seg> dq; // range: 1~C
  explicit DynamicHull(T _f, int _C): f(_f), C(_C) {} // max t s.t. f(x, t) >= f(y, t), x < y, maintain max
  int intersect(int x, int y) {
    int 1 = 0, r = C + 1;
    while (1 + 1 < r) {
      int mid = (1 + r) / 2;
      if (f(x, mid) >= f(y, mid)) l = mid;
      else r = mid;
    return 1:
  void push_back(int x) {
    for (int i; !dq.empty() &&
         (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
         ():
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back(). r + 1 \le C) dq.pb(seg({x, dq.back().}
         r + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
  }
};
     Flow & Matching
3
3.1 Dinic [801a71]
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
     \{2/3\}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
```

11 to, cap, flow, rev;

```
};
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
    n = _n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
    g[u].pb(edge{v, cap, 0, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
    dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int cur = q.front(); q.pop();
       for (auto &e : g[cur]) {
        if (dis[e.to] == -1 && e.flow != e.cap) {
           q.push(e.to);
           dis[e.to] = dis[cur] + 1;
      }
    }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
       edge &e = g[u][i];
       if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
         11 df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df;
           g[e.to][e.rev].flow -= df;
           return df;
         }
      }
    dis[u] = -1;
    return 0;
  11 maxflow(int _s, int _t) {
    s = _s; t = _t;
11 flow = 0, df;
    while (bfs()) {
       ind.assign(n, 0);
       while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
};
      Bounded Flow [758826]
3.2
struct BoundedFlow : Dinic {
  vector<ll> tot;
  void init(int _n) {
    Dinic::init(_n + 2);
    tot.assign(n, 0);
  void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
  bool feasible() {
    11 \text{ sum} = 0;
    int vs = n - 2, vt = n - 1;
for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
         add_edge(vs, i, 0, tot[i]), sum += tot[i];
       else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
```

dis.assign(n, INF);

```
if(tot[i] > 0)
                                                                 queue<int> q;
        g[vs].pop_back(), g[i].pop_back();
                                                                 auto relax = [&](int u, ll d, edge *e) {
      else if(tot[i] < 0)</pre>
                                                                   if (dis[u] > d) {
                                                                     dis[u] = d, past[u] = e;
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
                                                                     if (!inq[u]) inq[u] = 1, q.push(u);
                                                                   }
 11 boundedflow(int _s, int _t) {
  add_edge(_t, _s, 0, INF);
                                                                 };
                                                                 relax(s, 0, 0);
    if(!feasible()) return -1;
                                                                 while (!q.empty()) {
   ll x = g[_t].back().flow;
                                                                   int u = q.front();
                                                                   q.pop(), inq[u] = 0;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
                                                                   for (auto &e : g[u])
    //return x + maxflow(_s, _t); // max
                                                                     if (e.cap > e.flow)
                                                                        relax(e.to, dis[u] + e.cost, &e);
};
                                                                 }
3.3 MCMF [8d3644]
                                                               void try_edge(edge &cur) {
                                                                 if (cur.cap > cur.flow) return ++cur.cap, void();
struct MCMF { // O-based, O(SPFA * |f|)
                                                                 BellmanFord(cur.to);
  struct edge {
                                                                 if (dis[cur.from] + cur.cost < 0) {</pre>
    11 from, to, cap, flow, cost, rev;
                                                                   ++cur.flow, --g[cur.to][cur.rev].flow;
  };
                                                                   for (int i = cur.from; past[i]; i = past[i]->from
  int n;
  int s, t; ll mx;
                                                                     auto &e = *past[i];
 //mx: maximum amount of flow
                                                                      ++e.flow, --g[e.to][e.rev].flow;
  vector<vector<edge>> g;
                                                                   }
  vector<ll> dis, up;
                                                                 }
  bool BellmanFord(ll &flow, ll &cost) {
                                                                 ++cur.cap;
    vector<edge*> past(n);
                                                               }
    vector<int> inq(n);
                                                               void solve(int mxlg) { // mxlg >= log(max cap)
    dis.assign(n, INF); up.assign(n, 0);
                                                                 for (int b = mxlg; b >= 0; --b) {
    queue<int> q;
                                                                   for (int i = 0; i < n; ++i)</pre>
    q.push(s), inq[s] = 1;
                                                                     for (auto &e : g[i])
    up[s] = mx - flow, past[s] = 0, dis[s] = 0;
                                                                       e.cap *= 2, e.flow *= 2;
    while (!q.empty()) {
                                                                   for (int i = 0; i < n; ++i)</pre>
      int u = q.front();
                                                                     for (auto &e : g[i])
      q.pop(), inq[u] = 0;
                                                                       if (e.fcap >> b & 1)
      if (!up[u]) continue;
                                                                         try_edge(e);
      for (auto &e : g[u])
                                                                 }
        if (e.flow != e.cap &&
                                                               }
            dis[e.to] > dis[u] + e.cost) {
                                                               void init(int _n) {
          dis[e.to] = dis[u] + e.cost, past[e.to] = &e;
                                                                 n = _n;
          up[e.to] = min(up[u], e.cap
                                        - e.flow);
                                                                 past.assign(n, nullptr);
          if (!inq[e.to]) inq[e.to] = 1, q.push(e.to);
                                                                 g.assign(n, vector<edge>());
                                                               void add_edge(ll a, ll b, ll cap, ll cost) {
    if (dis[t] == INF) return 0;
                                                                 g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a)}
    flow += up[t], cost += up[t] * dis[t];
                                                                      == b)});
    for (ll i = t; past[i]; i = past[i] \rightarrow from) {
                                                                 g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1});
      auto &e = *past[i];
                                                               }
      e.flow += up[t], g[e.to][e.rev].flow -= up[t];
                                                             };
    }
    return 1;
                                                             3.5 Gomory Hu [82d968]
  pll MinCostMaxFlow(int _s, int _t) {
                                                             void GomoryHu(Dinic &flow) { // 0-based
    s = _s, t = _t;
11 flow = 0, cost = 0;
                                                               int n = flow.n;
                                                               vector<int> par(n);
    while (BellmanFord(flow, cost));
                                                               for (int i = 1; i < n; ++i) {</pre>
    return pll(flow, cost);
                                                                 flow.reset();
                                                                 add_edge(i, par[i], flow.maxflow(i, par[i]));
  void init(int _n, ll _mx) {
                                                                 for (int j = i + 1; j < n; ++j)
                                                                   if (par[j] == par[i] && ~flow.dis[j])
   n = n, mx = mx;
   g.assign(n, vector<edge>());
                                                                     par[j] = i;
 void add_edge(int a, int b, 11 cap, 11 cost) {
   g[a].pb(edge{a, b, cap, 0, cost, SZ(g[b])});
                                                             }
                                                             3.6 Stoer Wagner Algorithm [a9917b]
    g[b].pb(edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
                                                             struct StoerWagner { // 0-based, 0(V^3)
};
                                                               int n:
3.4 Min Cost Circulation [47cf18]
                                                               vector<int> vis, del;
                                                               vector<ll> wei;
struct MinCostCirculation { // 0-based, O(VE * ElogC)
                                                               vector<vector<ll>> edge;
  struct edge {
                                                               void init(int _n) {
   11 from, to, cap, fcap, flow, cost, rev;
                                                                 n = _n;
                                                                 del.assign(n, 0);
  int n;
                                                                 edge.assign(n, vector<ll>(n));
  vector<edge*> past;
  vector<vector<edge>> g;
                                                               void add_edge(int u, int v, ll w) {
  vector<ll> dis;
                                                                 edge[u][v] += w, edge[v][u] += w;
  void BellmanFord(int s) {
    vector<int> inq(n);
                                                               void search(int &s, int &t) {
```

vis.assign(n, 0); wei.assign(n, 0);

```
s = t = -1;
    while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)
        if (!del[i] && !vis[i] && mx < wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)</pre>
        if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
    }
  ll solve() {
    11 ret = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {</pre>
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
for (int j = 0; j < n; ++j)</pre>
        edge[x][j] = (edge[j][x] += edge[y][j]);
    return ret;
  }
};
3.7 Bipartite Matching [013c49]
//min vertex cover: take all unmatched vertices in L
    and find alternating tree,
//ans is not reached in L + reached in R
// O(VE)
int n; // 1-based, max matching
int mx[maxn], my[maxn];
bool adj[maxn][maxn], vis[maxn];
bool dfs(int u) {
  if (vis[u]) return 0;
  vis[u] = 1;
  for (int v = 1; v <= n; v++) {</pre>
    if (!adj[u][v]) continue;
    if (!my[v] || (my[v] && dfs(my[v]))) {
      mx[u] = v, my[v] = u;
      return 1:
    }
  }
  return 0;
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!\sim my[e] \mid | (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  }
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < nl; ++i)</pre>
      if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
        }
    }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
```

```
if (!~mx[i] && dfs(i)) ++ret;
    }
    return ret:
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
    nl = _nl, nr = _nr;
    g.assign(nl, vector<int>());
};
3.8
      Kuhn Munkres Algorithm [683e0a]
struct KM { // 0-based, maximum matching, O(V^3)
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll> hl, hr, slk;
  vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    n = _n;
// -INF for perfect matching
    w.assign(n, vector<ll>(n, 0));
    pre.assign(n, 0);
    qu.assign(n, 0);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
        0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
              w[x][y])) {
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !slk[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
        ; hr.assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
     hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
 }
};
3.9
     Max Simple Graph Matching [907d7c]
struct Matching { // 0-based, O(V^3)
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> g;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
```

vis[x] = tk;

```
x = Find(pre[match[x]]);
  }
void Blossom(int x, int y, int 1) {
  for (; Find(x) != 1; x = pre[y]) {
    pre[x] = y, y = match[x];
    if (s[y] == 1) q.push(y), s[y] = 0;
    for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
bool Bfs(int r) {
  iota(iter(fa), 0); fill(iter(s), -1);
  q = queue < int > (); q.push(r); s[r] = 0;
  for (; !q.empty(); q.pop()) {
    for (int x = q.front(); int u : g[x])
      if (s[u] == -1) {
        if (pre[u] = x, s[u] = 1, match[u] == n) {
           for (int a = u, b = x, last;
    b != n; a = last, b = pre[a])
             last = match[b], match[b] = a, match[a] =
                  b;
           return true;
        q.push(match[u]); s[match[u]] = 0;
      } else if (!s[u] && Find(u) != Find(x)) {
        int 1 = LCA(u, x);
        Blossom(x, u, 1); Blossom(u, x, 1);
  return false;
Matching(\textbf{int} \_n) : n(\_n), fa(n + 1), s(n + 1), vis(n
+ 1), pre(n + 1, n), match(n + 1, n), g(n) {} void add_edge(int u, int v)
{ g[u].pb(v), g[v].pb(u); }
int solve() {
  int ans = 0:
  for (int x = 0; x < n; ++x)
    if (match[x] == n) ans += Bfs(x);
  return ans:
} // match[x] == n means not matched
```

3.10 Stable Marriage

```
1: Initialize m \in M and w \in W to free
 2: while \exists free man m who has a woman w to propose to do
        w \leftarrow \mathsf{first} woman on m's list to whom m has not yet proposed
 4:
        if \exists some pair (m', w) then
 5:
            if w prefers m to m' then
 6:
                m' \leftarrow \textit{free}
7:
                 (m,w) \leftarrow \mathsf{engaged}
8:
            end if
9:
        else
10.
             (m,w) \leftarrow \mathsf{engaged}
         end if
11:
12: end while
```

3.11 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
 - 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect t o s with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow

- 1. Consruct super source S and sink T
- 2. For each edge (x,y,c), connect x o y with (cost,cap)=(c,1) if c>0, otherwise connect y o x with (cost,cap)=(-c,1)
- 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1. decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=(0,d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer ${\cal T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v$, $v \in \overset{\cdot}{G}$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u\to v$ and $v\to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. Let $w'(u,v)=w(u,v)-\mu(u)-\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 2. Find the minimum weight matching M with w'. The answer is $\sum \mu(v) + w'(M)$.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \geq 0$$

4 Geometry

4.1 Geometry Template [86f0f1]

```
using ld = 11;
using pdd = pair<ld, ld>;
#define X first
#define Y second
// Ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){ return v > eps ? 1 : ( v < -eps ? -1
    : 0);
int ori(pdd a, pdd b, pdd c)
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)
    ) <= 0; }
```

bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){

```
if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
      p2) || btw(p4, p1, p2))
    return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&
  ori(p3, p4, p1) * ori(p3, p4, p2) < 0;</pre>
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
    p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) /
      abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
      pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
       dot(dp, dq));
  return q0 + pdd(cross(r - p0, num), dot(r - p0, num))
        / abs2(dp);
} // from line p0--p1 to q0--q1, apply to r
```

4.2 Polar Angle Comparator [808e89]

4.3 Minkowski Sum [98abff]

```
void reorder_poly(vector<pdd>& pnts){
  int mn = 0;
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn</pre>
        ].Y && pnts[i].X < pnts[mn].X))</pre>
      mn = i:
 rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
 reorder_poly(P);
  reorder_poly(Q);
  int psz = P.size();
 int qsz = Q.size();
 P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
 vector<pdd> ans;
 int i = 0, j = 0;
 while(i < psz || j < qsz){
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1] - P[i], Q[j + 1] - Q[j])
    if(t >= 0) i++;
   if(t <= 0) j++;
 return ans;
}
```

4.4 Intersection of Circle and Convex Polygon [63653d]

```
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);
  if(abs(pb)<eps) return 0;
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
```

4.5 Intersection of Circles [f7a2fe]

4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
   int sign1 ){
  vector<Line> ret;
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
 pdd p1 = c1.0 + n * c1.R;
   pdd p2 = c2.0 + n * (c2.R * sign1);
   if (sgn(p1.X - p2.X) == 0 and
       sgn(p1.Y - p2.Y) == 0)
     p2 = p1 + perp(c2.0 - c1.0);
   ret.pb(Line(p1, p2));
  return ret:
}
```

4.7 Intersection of Line and Convex Polygon [157258]

```
int TangentDir(vector<pll> &C, pll dir) {
    return cyc_tsearch(SZ(C), [&](int a, int b) {
        return cross(dir, C[a]) > cross(dir, C[b]);
    });
}
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
    int A = TangentDir(C, a - b);
    int B = TangentDir(C, b - a);
    int n = SZ(C);
    if (cmpL(A) < 0 || cmpL(B) > 0)
        return pii(-1, -1); // no collision
    auto gao = [&](int l, int r) {
        for (int t = 1; (l + 1) % n != r; ) {
            int m = ((l + r + (l < r ? 0 : n)) / 2) % n;
            (cmpL(m) == cmpL(t) ? l : r) = m;
        }
        return (l + !cmpL(r)) % n;</pre>
```

```
};
pii res = pii(gao(B, A), gao(A, B)); // (i, j)
if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
    +1
    switch ((res.X - res.Y + n + 1) % n) {
        case 0: return pii(res.X, res.X);
        case 2: return pii(res.Y, res.Y);
}
/* crossing sides (i, i+1) and (j, j+1)
crossing corner i is treated as side (i, i+1)
returned in the same order as the line hits the
        convex */
return res;
} // convex cut: (r, l]
```

4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
    a, pdd b) {
  pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
  ;
  double s = cross(b - a, c - a), h2 = r * r - s * s /
    abs2(b - a);
  if (sgn(h2) < 0) return {};
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return {p - h, p + h};
}</pre>
```

4.9 Point in Circle [ecf954]

```
// return q's relation with circumcircle of tri(p[0],p
       [1],p[2])
bool in_cc(const array<pl1, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
       det += __int128(abs2(p[i]) - abs2(q)) * cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

4.10 Point in Convex [f86640]

4.11 Half Plane Intersection [dfb833]

```
// from 8BQube
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
     .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128) a02Y * a12X - (__int128) a02X *
       a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(iter(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
      return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
```

```
for (auto p : arr) {
 if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
       -1)
    continue;
  while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.
      back()))
    dq.pop_back();
  while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
    dq.pop_front();
 dq.pb(p);
while (SZ(dq) >= 3 \&\& !isin(dq[0], dq[SZ(dq) - 2], dq
    .back()))
  dq.pop_back();
while (SZ(dq) >= 3 && !isin(dq.back(), dq[0], dq[1]))
  dq.pop_front();
return vector<Line>(iter(dq));
```

4.12 Minimum Enclosing Circle [5af6d5]

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
  random_shuffle(iter(pts));
  pdd c = pts[0];
  ld r = 0;
  for(int i = 1; i < SZ(pts); i++){</pre>
    if(abs(pts[i] - c) <= r) continue;</pre>
    c = pts[i]; r = 0;
    for(int j = 0; j < i; j++){</pre>
      if(abs(pts[j] - c) <= r) continue;</pre>
      c = (pts[i] + pts[j]) / 2;
      r = abs(pts[i] - c);
      for(int k = 0; k < j; k++){</pre>
        if(abs(pts[k] - c) > r)
          tie(c, r) = circumcenter(pts[i], pts[j], pts[
               k]);
      }
   }
  return {c, r};
```

4.13 3D Point [badbbd]

.y * p.y), p.z); }

```
// Copy from 8BQube
struct Point {
  double x, y, z;
  Point(double \underline{x} = \emptyset, double \underline{y} = \emptyset, double \underline{z} = \emptyset): x = (\underline{x}), y(\underline{y}), z(\underline{z})
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
```

```
Point masscenter(Point a, Point b, Point c, Point d)
                                                             return fabs(ans / 6);
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
                                                           double get_dis(Point p, Face f) {
// proj. u to the plane of a, b, and c
 Point e1 = b - a;
                                                                  .z) * (p3.y - p1.y);
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
                                                                  .x) * (p3.z - p1.z);
  e2 = e2 / abs(e2);
 Point p = u - a;
 return pdd(dot(p, e1), dot(p, e2));
                                                                   * a + b * b + c * c);
Point rotate_around(Point p, double angle, Point axis)
  double s = sin(angle), c = cos(angle);
                                                           };
 Point u = axis / abs(axis);
  return u * dot(u, p) * (1 - c) + p * c + cross(u, p)
}
4.14 ConvexHull3D [156311]
                                                           /* Delaunay Triangulation:
struct convex_hull_3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
                                                           struct Edge {
}; // return the faces with pt indexes
                                                             int id; // oidx[id]
vector<Face> res;
                                                             list<Edge>::iterator twin;
                                                             Edge(int _id = 0):id(_id) {}
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
                                                           struct Delaunay { // 0-base
 int n = SZ(P);
                                                             int n:
                                                             vector<int> oidx;
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn
                                                             vector<pll> p;
      (abs2(P[0] - p)) != 0; }));
                                                                  head(n), p(n) {
  swap(P[2], *find_if(iter(P), [&](auto p) { return sgn
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
                                                                iota(iter(oidx), 0);
  swap(P[3], *find_if(iter(P), [&](auto p) { return sgn
      (volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
                                                                divide(0, n - 1);
    vector<Face> next;
    for (auto f : res) {
                                                             void addEdge(int u, int v) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
                                                               head[u].push_front(Edge(v));
                                                                head[v].push_front(Edge(u));
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
                                                             void divide(int 1, int r) {
                                                                if (1 == r) return;
    for (auto f : res) {
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
          next.emplace_back(x, y, i);
                                                                auto gao = [&](int t) {
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
    res = next;
 }
bool same(Face s, Face t) {
 if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
                                                                  return false:
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
                                                                while (gao(0) || gao(1));
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
                                                                while (true) {
       return 0;
                                                                  int ch = -1, sd = 0;
  return 1;
                                                                  for (int t = 0; t < 2; ++t)
int polygon_face_num() {
 int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin() + i, [&](
                                                                          id])))
       Face g) { return same(res[i], g); });
                                                                        ch = it.id, sd = t;
  return ans;
double get_volume() {
                                                                      nw[sd]].end(); )
 double ans = 0;
                                                                        sd ^ 1], p[ch]))
  for (auto f : res)
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
        ]);
```

```
Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
  double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
  .y) * (p3.x - p1.x);
double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
// n^2 delaunay: facets with negative z normal of // convexhull of (x, y, x^2 + y^2), use a pseudo-point
// (0, 0, inf) to avoid degenerate case
4.15 Delaunay Triangulation [6a9916]
   Given a sets of points on 2D plane, find a
   triangulation such that no points will strictly
   inside circumcircle of any triangle. */
  vector<list<Edge>> head; // result udir. graph
  Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
    for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(iter(oidx), [&](int a, int b)
          return _p[a] < _p[b]; });</pre>
    for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
    if (l + 1 == r) return addEdge(l, l + 1);
    int mid = (1 + r) >> 1, nw[2] = \{1, r\};
    divide(l, mid), divide(mid + 1, r);
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      for (auto it : head[nw[t]]) {
        int v = ori(pt[1], pt[0], p[it.id]);
        if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
             id]) < abs2(pt[1] - pt[0])))
          return nw[t] = it.id, true;
    addEdge(nw[0], nw[1]); // add tangent
      pll pt[2] = {p[nw[0]], p[nw[1]]};
        for (auto it : head[nw[t]])
          if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
                -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
      if (ch == -1) break; // upper common tangent
      for (auto it = head[nw[sd]].begin(); it != head[
        if (seg_strict_intersect(pt[sd], p[it->id], pt[
           head[it->id].erase(it->twin), head[nw[sd]].
               erase(it++);
```

```
double m = 2. * (x1 * y2 - y1 * x2);
      nw[sd] = ch, addEdge(nw[0], nw[1]);
                                                              pdd center:
                                                              center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
    }
                                                              y1 - y2)) / m;
center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
 }
};
                                                                   y2 * y2) / m;
4.16 Voronoi Diagram [e4f408]
                                                              return center + p0;
// all coord. is even, you may want to call
                                                            pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
    halfPlaneInter after then
vector<vector<Line>> vec;
                                                              double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
void build_voronoi_line(int n, vector<pll> &pts) {
                                                                    - p2);
  Delaunay tool(n, pts); // Delaunay
                                                              double s = a + b + c;
  vec.clear(), vec.resize(n);
                                                              return (a * p1 + b * p2 + c * p3) / s;
  for (int i = 0; i < n; ++i)
    for (auto e : tool.head[i]) {
                                                            pdd masscenter(pdd p1, pdd p2, pdd p3)
      int u = tool.oidx[i], v = tool.oidx[e.id];
                                                            { return (p1 + p2 + p3) / 3; }
      pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v]
                                                            pdd orthcenter(pdd p1, pdd p2, pdd p3)
           - pts[u]);
                                                            { return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
      vec[u].pb(Line(m, m + d));
}
                                                            4.20 Rotating Sweep Line [f5f689]
4.17 Polygon Union [9fbf66]
                                                            struct Event {
                                                              pll d; int u, v;
// from 8BOube
                                                              bool operator<(const Event &b) const {</pre>
ld rat(pll a, pll b) {
                                                                int ret = cmp(d, b.d, false);
 return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
                                                                return ret == -1 ? false : ret; } // no tie-break
} // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
                                                            void rotatingSweepLine(const vector<pll> &p) {
  ld res = 0;
                                                              const int n = SZ(p);
  for (auto &p : poly)
                                                              vector<Event> e; e.reserve(n * (n - 1));
    for (int a = 0; a < SZ(p); ++a) {</pre>
                                                              for (int i = 0; i < n; i++)</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
                                                                for (int j = 0; j < n; j++) // pos[i] < pos[j] when</pre>
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
                                                                      the event occurs
      for (auto &q : poly) {
                                                                  if (i != j) e.pb(p[j] - p[i], i, j);
        if (&p == &q) continue;
                                                              sort(iter(e));
        for (int b = 0; b < SZ(q); ++b) {</pre>
                                                              vector<int> ord(n), pos(n);
          pll C = q[b], D = q[(b + 1) \% SZ(q)];
                                                              iota(iter(ord), 0);
          int sc = ori(A, B, C), sd = ori(A, B, D);
                                                              sort(iter(ord), [&](int i, int j) { // initial order
          if (sc != sd && min(sc, sd) < 0) {
                                                                   return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
            1d sa = cross(D - C, A - C), sb = cross(D - C)
                                                                       X < p[j].X; \});
                 C, B - C);
                                                              for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
            segs.pb(sa / (sa - sb), sgn(sc - sd));
                                                              // initialize
                                                              for (int i = 0, j = 0; i < SZ(e); i = j) {
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D</pre>
                                                                // do somethina
                - C)) > 0) {
                                                                vector<pii> tmp;
            segs.pb(rat(C - A, B - A), 1);
                                                                for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
            segs.pb(rat(D - A, B - A), -1);
                                                                  tmp.pb(pii(e[j].u, e[j].v));
                                                                 sort(iter(tmp), [&](pii x, pii y){
        }
                                                                    return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
      }
                                                                         ], pos[y.ss]); });
      sort(iter(segs));
                                                                for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
  tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
                                                                    make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
      int cnt = segs[0].second;
                                                                         pos[x]);
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
                                                            }
        cnt += segs[j].Y;
                                                            4.21 Vector In Poly [c6d0fa]
      res += cross(A, B) * sum;
                                                            // ori(a, b, c) >= 0, valid: "strict" angle from a-b to
  return res / 2;
                                                                  a-c
                                                            bool btwangle(pll a, pll b, pll c, pll p, int strict) {
                                                              return ori(a, b, p) >= strict && ori(a, p, c) >=
4.18 Tangent Point to Convex Hull [523bc1]
                                                                   strict;
// from 8BOube
                                                            // whether vector{cur, p} in counter-clockwise order
/* The point should be strictly out of hull
                                                                prv, cur, nxt
 return arbitrary point on the tangent line */
                                                            bool inside(pll prv, pll cur, pll nxt, pll p, int
pii get_tangent(vector<pll> &C, pll p) {
                                                                 strict) {
  auto gao = [&](int s) {
                                                              if (ori(cur, nxt, prv) >= 0)
    return cyc_tsearch(SZ(C), [&](int x, int y)
                                                                return btwangle(cur, nxt, prv, p, strict);
    { return ori(p, C[x], C[y]) == s; });
                                                              return !btwangle(cur, prv, nxt, p, !strict);
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
                                                            4.22 Convex Hull DP [92fd4b]
4.19 Heart [082d19]
                                                            sort(iter(pts), [&](pll x, pll y) {
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
                                                                return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;</pre>
    (center)
                                                            auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
```

]; };

```
// DP for convex hull vertices (no points on edges)
auto solve = [\&](int bottom) { // <math>O(n^3)
  pll 0 = pts[bottom];
  vector<pii> trans;
  for (int j = bottom + 1; j < n; j++)</pre>
    for (int k = bottom + 1; k < n; k++) {
      if (ori(0, pts[j], pts[k]) <= 0) continue;</pre>
      // check whether j->k is legal
      trans.pb(pii(j, k));
  sort(iter(trans), [&](pii x, pii y) -> bool{
      int tmp = cmp(getvec(x), getvec(y), false);
      if (tmp != -1) return tmp;
      pll v = getvec(x);
      return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether bottom -> j is legal
    // init trans -> j
  for (auto [i, j] : trans) {
    // normal trans i -> j
  for (int j = bottom + 1; j < n; j++) {</pre>
   // check whether j -> bottom is legal
    // end trans j ->
 }
};
for(int i = 0; i < n; i++) solve(i);</pre>
4.23 Calculate Points in Triangle [bf746f]
// all points are distinct
// cnt[i][j] = # of point k s.t. strictly above ij, and
     i < k < j
// cnt2[i][j] = # of points k s.t. strictly in ij
// preprocess space: O(n^2), time: O(n^3), query time:
    0(1)
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
  for (int j = 0; j < n; j++){
    if (pts[i] >= pts[j]) continue;
    for (int k = 0; k < n; k++) {
      if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
        int tmp = ori(pts[i], pts[j], pts[k]);
if (tmp > 0) cnt[i][j]++; // only for i < j</pre>
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
      }
   }
  }
auto calc_tri = [&](array<int, 3> arr) { // strictly
    inside
  sort(iter(arr), [&](int x, int y){ return pts[x] <</pre>
      pts[y]; });
  auto [x, y, z] = arr;
  int tmp = ori(pts[x], pts[y], pts[z]);
  if (tmp == 0) return 0;
  else if (tmp < 0)</pre>
    y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
      x][z];
};
5
     Graph
      BCC [d04ebe]
struct BCC{ // O-based, allow multi edges but not allow
     Loops
  int n, m, cnt = 0;
  // n:|V|, m:|E|, cnt:#bcc
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
 vector<vector<int>> bcc_v, bcc_e;
vector<vector<pii>> g; // original graph
  vector<pii> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
```

n(_n), m(SZ(_edges)), g(_n), edges(_edges){

g[u].pb(pii(v, i)); g[v].pb(pii(u, i));

for(int i = 0; i < m; i++){</pre>

auto [u, v] = edges[i];

```
void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
  // modify these if you need more information
  void add_v(int v){ bcc_v.back().pb(v); }
  void add_e(int e){ bcc_e.back().pb(e); }
  void build(){
    vector<int> in(n, -1), low(n, -1), stk;
    vector<vector<int>> up(n);
    int ts = 0;
    auto _dfs = [&](auto dfs, int now, int par, int pe)
         -> void{
      if(pe != -1) up[now].pb(pe);
      in[now] = low[now] = ts++;
      stk.pb(now);
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
        if(in[v] != -1){
          if(in[v] < in[now]) up[now].pb(e);</pre>
          low[now] = min(low[now], in[v]);
          continue:
        dfs(dfs, v, now, e);
        low[now] = min(low[now], low[v]);
      if((now != par && low[now] >= in[par]) || (now ==
           par && SZ(g[now]) == 0)){
        make_bcc();
        for(int v = stk.back();; v = stk.back()){
          stk.pop_back(), add_v(v);
          for(int e : up[v]) add_e(e);
          if(v == now) break;
        if(now != par) add_v(par);
      }
    };
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) _dfs(_dfs, i, i, -1);
  }
};
5.2 SCC [2c9a01]
struct SCC{ // 0-based, output reversed topo order
  int n, cnt = 0;
  vector<vector<int>> g;
  vector<int> sccid;
  explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
  void add_edge(int u, int v){
    g[u].pb(v);
  void build(){
    vector<int> in(n, -1), low(n), stk;
    vector<bool> instk(n);
    int ts = 0;
    auto dfs1 = [&](auto dfs, int now) -> void{
      stk.pb(now); instk[now] = true;
      in[now] = low[now] = ts++;
      for(int i : g[now]){
        if(in[i] == -1)
          dfs(dfs, i), low[now] = min(low[now], low[i])
        else if(instk[i] && in[i] < in[now])</pre>
          low[now] = min(low[now], in[i]);
      if(low[now] == in[now]){
        for(; stk.back() != now; stk.pop_back())
          sccid[stk.back()] = cnt, instk[stk.back()] =
               false;
        sccid[now] = cnt++, instk[now] = false, stk.
            pop_back();
      }
    };
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) dfs1(dfs1, i);
};
5.3
     2-SAT [0686a5]
struct SAT { // 0-based
  int n;
  vector<bool> istrue;
  SCC scc;
```

void insert(int u) {

int p = LCA(st[top], u);

if(dep[vrt] > dep[p]) vrt = p;

if (top == -1) return st[++top] = vrt = u, void();

if (p == st[top]) return st[++top] = u, void();

```
_n): n(_n), istrue(n + n), scc(n + n) {}
                                                                while (top >= 1 \&\& dep[st[top - 1]] >= dep[p])
                                                                  vG[st[top - 1]].pb(st[top]), --top;
  int neg(int a) {
    return a >= n ? a - n : a + n;
                                                                if (st[top] != p)
                                                                  vG[p].pb(st[top]), --top, st[++top] = p;
  void add_clause(int a, int b) {
                                                                st[++top] = u;
    scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
  bool solve() {
                                                              void reset(int u) {
                                                                for (int i : vG[u]) reset(i);
    scc.build();
    for (int i = 0; i < n; ++i) {</pre>
                                                                vG[u].clear();
      if (scc.sccid[i] == scc.sccid[i + n]) return
          false;
      istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>
                                                             void solve(vector<int> &v) {
      istrue[i + n] = !istrue[i];
                                                                top = -1;
                                                                sort(ALL(v),
                                                                    [&](int a, int b) { return dfn[a] < dfn[b]; });
    return true;
                                                                for (int i : v) insert(i);
};
                                                                while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
                                                                // do something
5.4 Dominator Tree [2da9bb]
                                                                reset(vrt);
                                                             }
struct Dominator {
                                                             5.6 Fast DMST [7b274d]
  int n;
  vector<vector<int>> g, r, rdom; int tk;
  vector<int> dfn, rev, fa, sdom, dom, val, rp;
                                                              struct E { int s, t; ll w; }; // O-base
  Dominator(int _n): n(_n), g(n), r(n), rdom(n), tk(0)
                                                              struct PQ {
                                                                struct P
    dfn = rev = fa = sdom = dom =
                                                                  11 v; int i;
      val = rp = vector<int>(n, -1); }
                                                                  bool operator>(const P &b) const { return v > b.v;
  void add_edge(int x, int y) { g[x].push_back(y); }
  void dfs(int x) {
    rev[dfn[x] = tk] = x;
                                                                priority_queue<P, vector<P>, greater<>> pq; 11 tag;
    fa[tk] = sdom[tk] = val[tk] = tk; tk++;
                                                                    // min heap
                                                                void push(P p) { p.v -= tag; pq.emplace(p); }
    for (int u : g[x]) {
      if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                                P top() { P p = pq.top(); p.v += tag; return p; }
                                                                void join(PQ &b) {
      r[dfn[u]].push_back(dfn[x]);
    }
                                                                  if (pq.size() < b.pq.size())</pre>
                                                                    swap(pq, b.pq), swap(tag, b.tag);
  void merge(int x, int y) { fa[x] = y; }
                                                                  while (!b.pq.empty()) push(b.top()), b.pq.pop();
  int find(int x, int c = 0) {
    if (fa[x] == x) return c ? -1 : x;
                                                             }; // O(E log^2 V), use leftist tree for O(E log V)
    if (int p = find(fa[x], 1); p != -1) {
                                                             vector<int> dmst(const vector<E> &e, int n, int root) {
                                                               vector<PQ> h(n * 2);
for (int i = 0; i < int(e.size()); ++i)</pre>
      if (sdom[val[x]] > sdom[val[fa[x]]])
        val[x] = val[fa[x]];
      fa[x] = p;
                                                                  h[e[i].t].push({e[i].w, i});
                                                               vector<int> a(n * 2); iota(iter(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
      return c ? p : val[x];
    } else return c ? fa[x] : val[x];
                                                                auto o = [\&](auto Y, int x) \rightarrow int {
                                                                  return x == a[x] ? x : a[x] = Y(Y, a[x]); };
  vector<int> build(int s) {
    // return the father of each node in dominator tree dfs(s); // p[i] = -2 if i is unreachable, par[s] =
                                                                auto S = [&](int i) { return o(o, e[i].s); };
                                                                int pc = v[root] = n;
                                                                for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
    for (int i = tk - 1; i >= 0; --i) {
                                                                  for (int p = i; v[p]<0 \mid \mid v[p]==i; p = S(r[p])) {
      for (int u : r[i])
                                                                    if (v[p] == i)
        sdom[i] = min(sdom[i], sdom[find(u)]);
                                                                      for (int q = pc++; p != q; p = S(r[p])) {
      if (i) rdom[sdom[i]].push_back(i);
                                                                        h[p].tag -= h[p].top().v; h[q].join(h[p]);
      for (int u : rdom[i]) {
                                                                        pa[p] = a[p] = q;
        int p = find(u);
        dom[u] = (sdom[p] == i ? i : p);
                                                                    while (S(h[p].top().i) == p) h[p].pq.pop();
                                                                    v[p] = i; r[p] = h[p].top().i;
      if (i) merge(i, rp[i]);
                                                                vector<int> ans;
                                                                for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
    vector < int > p(n, -2); p[s] = -1;
    for (int i = 1; i < tk; ++i)</pre>
                                                                  for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                                      1)
    for (int i = 1; i < tk; ++i)</pre>
                                                                    v[f] = n;
      p[rev[i]] = rev[dom[i]];
                                                                  ans.push_back(r[i]);
    return p;
                                                                return ans; // default minimize, returns edgeid array
};
                                                             }
      Virtual Tree [f977f8]
                                                             5.7 Vizing [f4ebad]
// copy from 8BQube
                                                              // find D+1 edge coloring of a graph with max deg D
                                                             struct vizing { // returns edge coloring in adjacent
vector<int> vG[N];
                                                                  matrix G. 1 - based
int top, st[N];
                                                                const int N = 105;
int vrt = -1;
```

int C[N][N], G[N][N], X[N], vst[N], n; // ans: G[i][j

void init(int _n) { n = _n; // n = |V|+1

for (int i = 0; i <= n; ++i)</pre>

for (int j = 0; j <= n; ++j)</pre>

C[i][j] = G[i][j] = 0;

]._Find_next(p))

dfs(r, c, l + 1, mask);

r[tp] = p, c[tp] = k, ++tp;

```
void dfs(vector<int> &r, vector<int> &c, int 1,
  void solve(vector<pii> &E) {
                                                                   bitset<N> mask) {
    auto update = [&](int u)
                                                                 while (!r.empty()) {
    { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
                                                                   int p = r.back();
    auto color = [&](int u, int v, int c) {
                                                                   r.pop_back(), mask[p] = 0;
      int p = G[u][v];
                                                                   if (q + c.back() <= ans) return;</pre>
      G[u][v] = G[v][u] = c;
                                                                   cur[q++] = p;
      C[u][c] = v, C[v][c] = u;
                                                                   vector<int> nr;
      C[u][p] = C[v][p] = 0;
                                                                   for (int i : r) if (G[p][i]) nr.pb(i);
      if (p) X[u] = X[v] = p;
                                                                   if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else update(u), update(v);
                                                                   else if (q > ans) ans = q, copy_n(cur, q, sol);
      return p;
                                                                   c.pop_back(), --q;
                                                                 }
    };
    auto flip = [&](int u, int c1, int c2) {
      int p = C[u][c1];
                                                               int solve() {
      swap(C[u][c1], C[u][c2]);
                                                                 vector<int> r(n);
      if (p) G[u][p] = G[p][u] = c2;
                                                                 ans = q = 0, iota(ALL(r), 0);
      if (!C[u][c1]) X[u] = c1;
                                                                 pre_dfs(r, 0, bitset<N>(string(n, '1')));
      if (!C[u][c2]) X[u] = c2;
                                                                 return ans;
      return p;
                                                              }
                                                            };
    }:
    fill_n(X + 1, n, 1);
                                                            5.9
                                                                   Number of Maximal Clique [11fa26]
    for (int t = 0; t < SZ(E); ++t) {</pre>
      int u = E[t].X, v0 = E[t].Y, v = v0, c0 = X[u], c
                                                             struct BronKerbosch { // 1-base
           = c0, d;
                                                               int n, a[N], g[N][N];
      vector<pii> L;
                                                               int S, all[N][N], some[N][N], none[N][N];
      fill_n(vst + 1, n, 0);
                                                               void init(int _n) {
      while (!G[u][v0]) {
                                                                 n = _n;
for (int i = 1; i <= n; ++i)</pre>
        L.emplace_back(v, d = X[v]);
        if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
                                                                   for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
             --a) c = color(u, L[a].X, c);
        else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
                                                               void add_edge(int u, int v) {
             0; --a) color(u, L[a].X, L[a].Y);
                                                                 g[u][v] = g[v][u] = 1;
        else if (vst[d]) break;
        else vst[d] = 1, v = C[u][d];
                                                               void dfs(int d, int an, int sn, int nn) {
                                                                 if (S > 1000) return; // pruning
      if (!G[u][v0]) {
                                                                 if (sn == 0 && nn == 0) ++S;
        for (; v; v = flip(v, c, d), swap(c, d));
                                                                 int u = some[d][0];
        if (int a; C[u][c0]) {
                                                                 for (int i = 0; i < sn; ++i) {</pre>
          for (a = SZ(L) - 2; a >= 0 && L[a].Y != c; --
                                                                   int v = some[d][i];
                                                                   if (g[u][v]) continue;
          for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
                                                                   int tsn = 0, tnn = 0;
                                                                   copy_n(all[d], an, all[d + 1]);
        else --t;
                                                                   all[d + 1][an] = v;
     }
                                                                   for (int j = 0; j < sn; ++j)
   }
                                                                     if (g[v][some[d][j]])
 }
                                                                       some[d + 1][tsn++] = some[d][j];
};
                                                                   for (int j = 0; j < nn; ++j)</pre>
5.8
     Maximum Clique [d50aa9]
                                                                     if (g[v][none[d][j]])
                                                                       none[d + 1][tnn++] = none[d][j];
struct MaxClique { // fast when N <= 100</pre>
                                                                   dfs(d + 1, an + 1, tsn, tnn);
 bitset<N> G[N], cs[N];
                                                                   some[d][i] = 0, none[d][nn++] = v;
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
                                                               }
   n = _n;
for (int i = 0; i < n; ++i) G[i].reset();</pre>
                                                               int solve() {
                                                                 iota(some[0], some[0] + n, 1);
                                                                 S = 0, dfs(0, 0, n, 0);
  void add_edge(int u, int v) {
                                                                 return S;
   G[u][v] = G[v][u] = 1;
                                                               }
                                                            };
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
                                                            5.10 Minimum Mean Cycle [3e5d2b]
      for (int i : r) d[i] = (G[i] & mask).count();
                                                             // from 8BQube
      sort(ALL(r), [\&](int x, int y) \{ return d[x] > d[
                                                            11 road[N][N]; // input here
          y]; });
                                                            struct MinimumMeanCycle {
    vector<int> c(SZ(r));
                                                               11 dp[N + 5][N], n;
                                                               pll solve() {
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
                                                                 11 a = -1, b = -1, L = n + 1;
    cs[1].reset(), cs[2].reset();
                                                                 for (int i = 2; i <= L; ++i)</pre>
    for (int p : r) {
      int k = 1;
                                                                   for (int k = 0; k < n; ++k)
      while ((cs[k] & G[p]).any()) ++k;
                                                                     for (int j = 0; j < n; ++j)</pre>
      if (k > rgt) cs[++rgt + 1].reset();
                                                                       dp[i][j] =
      cs[k][p] = 1;
                                                                         min(dp[i - 1][k] + road[k][j], dp[i][j]);
                                                                 for (int i = 0; i < n; ++i) {</pre>
      if (k < 1ft) r[tp++] = p;
                                                                   if (dp[L][i] >= INF) continue;
                                                                   11 ta = 0, tb = 1;
for (int j = 1; j < n; ++j)</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k]</pre>
```

if (dp[j][i] < INF &&</pre>

if (ta == 0) continue;

ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)

ta = dp[L][i] - dp[j][i], tb = L - j;

```
if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      11 g = 1
               _gcd(a, b);
      return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
   n = _n;
for (int i = 0; i < n; ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
  }
};
5.11 Minimum Steiner Tree [21acea]
// from 8BQube
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
    static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcost[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
    }
  }
  void add_edge(int ui, int vi, int wi) {
    dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
           dst[i][j] =
             min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
           dp[msk][i] =
             vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
              submsk = (submsk - 1) \& msk)
           dp[msk][i] = min(dp[msk][i],
             dp[submsk][i] + dp[msk ^ submsk][i] -
               vcost[i]);
      for (int i = 0; i < n; ++i) {</pre>
        tdst[i] = INF;
         for (int j = 0; j < n; ++j)
          tdst[i] =
             min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    int ans = INF;
    for (int i = 0; i < n; ++i)
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans:
};
5.12 Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
```

```
for (int y : D[x]) vis[y] = 1;
for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
for (int y : D[x]) vis[y] = 0;
```

```
for (int x : ord) { // c4
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) c4 += vis[z]++;
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M))
```

// ax+ny = 1, $ax+ny == ax == 1 \ (mod \ n)$

if (y == 0) g=x,a=1,b=0;

void extgcd(ll x,ll y,ll &g,ll &a,ll &b) {

Math 6

Extended Euclidean Algorithm [c51ae9] 6.1

```
else extgcd(y,x%y,g,b,a),b-=(x/y)*a;
}
6.2 Floor & Ceil [134881]
ll ifloor(ll a,ll b){
  return a / b - (a % b && (a < 0) ^ (b < 0));
11 iceil(11 a,11 b){
  return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre [4e4b23]
```

```
Legendre symbol,
  such that the bottom doesn't need to be prime.
// (n|p) -> same as Legendre
// (n|ab) = (n|a)(n|b)
// work with Long Long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a >>= r;
```

// the Jacobi symbol is a generalization of the

```
if (a \& m \& 2) s = -s;
    swap(a, m);
  return s;
// 0: a == 0
// -1: a isn't a quad res of p
// else: return X with X^2 % p == a
// doesn't work with long long
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
  int b, d;
  for (; ; ) {
    b = rand() % p;
    d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
       )) % p;
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
```

6.4 Simplex [aa7741]

return g0;

```
// maximize c^T x
// subject to Ax <= b, x >= 0
// and stores the solution;
typedef long double T; // Long double, Rational, double
     + mod<P>...
```

int DiscreteLog(int s, int x, int y, int m) {

```
typedef vector<T> vd;
                                                                   constexpr int kStep = 32000;
typedef vector<vd> vvd;
                                                                   unordered_map<int, int> p;
                                                                   int b = 1;
const T eps = 1e-9, inf = 1/.0;
                                                                   for (int i = 0; i < kStep; ++i) {</pre>
#define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
                                                                     p[y] = i;
    y = 1LL * y * x % m;
                                                                     b = 1LL * b * x % m;
#define rep(i, l, n) for(int i = l; i < n; i++)
                                                                   for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
struct LPSolver {
  int m, n;
                                                                     if (p.find(s) != p.end()) return i + kStep - p[s];
  vector<int> N, B;
  vvd D;
                                                                   return -1:
  LPSolver(const vvd& A, const vd& b, const vd& c) :
    m(SZ(b)), n(SZ(c)), N(n+1), B(m), D(m+2, vd(n+2)) {
                                                                 int DiscreteLog(int x, int y, int m) {
       rep(i,0,m) \ rep(j,0,n) \ D[i][j] = A[i][j];
                                                                   if (m == 1) return 0;
       rep(i,0,m) { B[i] = n+i; D[i][n] = -1; D[i][n+1]
                                                                   int s = 1;
           = b[i]:}
                                                                   for (int i = 0; i < 100; ++i) {
       rep(j,0,n) \{ N[j] = j; D[m][j] = -c[j]; \}
                                                                     if (s == y) return i;
                                                                     s = 1LL * s * x % m;
       N[n] = -1; D[m+1][n] = 1;
                                                                   if (s == y) return 100;
                                                                   int p = 100 + DiscreteLog(s, x, y, m);
  void pivot(int r, int s) {
    T *a = D[r].data(), inv = 1 / a[s];
                                                                   if (fpow(x, p, m) != y) return -1;
    rep(i,0,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
                                                                   return p; //returns: x^p = y \pmod{m}
       T *b = D[i].data(), inv2 = b[s] *
       rep(j,0,n+2) b[j] -= a[j] * inv2;
                                                                 6.7 Miller Rabin & Pollard Rho [d3ecd2]
       b[s] = a[s] * inv2;
    }
                                                                 // n < 4,759,123,141 3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
    rep(j,0,n+2) if (j != s) D[r][j] *= inv;
    rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
                                                                 // n < 3,474,749,660,383 6 : primes <= 13
    D[r][s] = inv;
                                                                 // n < 2^64
     swap(B[r], N[s]);
                                                                 // 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                 11 mul(ll a, ll b, ll n){
                                                                   return (__int128)a * b % n;
  bool simplex(int phase) {
    int x = m + phase - 1;
                                                                 bool Miller_Rabin(ll a, ll n) {
     for (;;) {
                                                                   if ((a = a % n) == 0) return 1;
       int s = -1;
                                                                   if (n % 2 == 0) return n == 2;
       rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
                                                                   11 \text{ tmp} = (n - 1) / ((n - 1) & (1 - n));
       if (D[x][s] >= -eps) return true;
                                                                   ll t = __lg(((n - 1) & (1 - n))), x = 1;
       int r = -1;
                                                                   for (; tmp; tmp >>= 1, a = mul(a, a, n))
       rep(i,0,m) {
                                                                     if (tmp & 1) x = mul(x, a, n);
         if (D[i][s] <= eps) continue;</pre>
                                                                   if (x == 1 || x == n - 1) return 1;
         if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])
                                                                   while (--t)
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
                                                                     if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                   return 0;
       if (r == -1) return false;
       pivot(r, s);
                                                                 bool prime(ll n){
    }
                                                                   vector<ll> tmp = {2, 325, 9375, 28178, 450775,
  }
                                                                        9780504, 1795265022};
                                                                   for(ll i : tmp)
  T solve(vd &x) {
                                                                     if(!Miller_Rabin(i, n)) return false;
    int r = 0;
                                                                   return true;
    rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] < -eps) {</pre>
                                                                 map<ll, int> cnt;
       pivot(r, n);
                                                                 void PollardRho(ll n) {
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                   if (n == 1) return;
           inf;
                                                                   if (prime(n)) return ++cnt[n], void();
       rep(i,0,m) if (B[i] == -1) {
                                                                   if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
                                                                        void();
         rep(j,1,n+1) ltj(D[i]);
                                                                   11 x = 2, y = 2, d = 1, p = 1;
         pivot(i, s);
                                                                 #define f(x, n, p) ((mul(x, x, n) + p) % n)
      }
                                                                   while (true) {
  if (d != n && d != 1) {
    bool ok = simplex(1); x = vd(n);
                                                                        PollardRho(n / d);
    rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
                                                                        PollardRho(d);
    return ok ? D[m][n+1] : inf;
                                                                        return;
  }
};
                                                                     if (d == n) ++p;
                                                                     x = f(x, n, p), y = f(f(y, n, p), n, p);
d = gcd(abs(x - y), n);
6.5
      Simplex Construction
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for
                                                                   }
all 1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
                                                                 }
1. In case of minimization, let c_i' = -c_i
                                                                 6.8 XOR Basis [006505]
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j \to \mathsf{add} \subseteq \mathsf{and} \supseteq.
                                                                 const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                 struct Basis{
6.6 DiscreteLog [da27bf]
                                                                   int total = 0, rank = 0;
                                                                   vector<11> b;
```

Basis(): b(digit) {}

```
bool add(ll v){ // Gauss Jordan Elimination
    total++:
    for(int i = digit - 1; i >= 0; i--){
       if(!(1LL << i & v)) continue;</pre>
      if(b[i] != 0){
         v ^= b[i];
         continue;
       for(int j = 0; j < i; j++)
       if(1LL << j & v) v ^= b[j];
for(int j = i + 1; j < digit; j++)</pre>
         if(1LL << i & b[j]) b[j] ^= v;</pre>
      b[i] = v;
      rank++;
      return true;
    }
    return false;
  11 \text{ getmax}(11 \text{ x} = 0){
    for(ll i : b) x = max(x, x ^ i);
    return x;
  11 \text{ getmin}(11 \text{ x} = 0)
    for(ll i : b) x = min(x, x ^ i);
    return x;
  bool can(ll x){
    return getmin(x) == 0;
  11 kth(11 k){ // kth smallest, 0-indexed
    vector<ll> tmp;
    for(ll i : b) if(i) tmp.pb(i);
    11 \text{ ans} = 0;
    for(int i = 0; i < SZ(tmp); i++)</pre>
      if(1LL << i & k) ans ^= tmp[i];</pre>
    return ans;
  }
};
6.9 Linear Equation [ab487b]
  int N = mat.size(), M = mat[0].size();
  int rk = 0;
```

```
vector<int> RREF(vector<vector<ll>> &mat){
  vector<int> cols;
  for (int i = 0;i < M;i++) {</pre>
    int cnt = -1;
    for (int j = N-1; j >= rk; j--)
      if(mat[j][i] != 0) cnt = j;
    if(cnt == -1) continue;
    swap(mat[rk], mat[cnt]);
    11 lead = mat[rk][i];
    for (int j = 0;j < M;j++) mat[rk][j] = mat[rk][j] *</pre>
         modinv(lead) % mod;
    for (int j = 0; j < N; j++) {</pre>
      if(j == rk) continue;
      11 tmp = mat[j][i];
      for (int k = 0; k < M; k++)
        mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % mod
              + mod) % mod;
    }
    cols.pb(i);
   rk++;
  }
  return cols;
struct LinearEquation{
  bool ok;
  vector<ll> par; //particular solution (Ax = b)
  vector<vector<ll>> homo; //homogenous (Ax = 0)
  vector<vector<ll>> rref;
  //first M columns are matrix A
  //last column of eq is vector b
  void solve(const vector<vector<ll>>> &eq){
    int M = (int)eq[0].size() - 1;
    rref = eq;
    auto piv = RREF(rref);
    int rk = piv.size();
    if(piv.size() && piv.back() == M){
      ok = 0; return;
    }
    ok = 1;
    par.resize(M);
```

```
vector<bool> ispiv(M);
    for (int i = 0;i < rk;i++) {</pre>
      par[piv[i]] = rref[i][M];
      ispiv[piv[i]] = 1;
    for (int i = 0;i < M;i++) {</pre>
      if (ispiv[i]) continue;
      vector<ll> h(M);
      h[i] = 1;
      for (int j = 0;j < rk;j++) h[piv[j]] = rref[j][i]</pre>
            ? mod-rref[j][i] : 0;
      homo.pb(h);
  }
};
```

6.10 Chinese Remainder Theorem [6ef4a3]

```
pll solve_crt(ll x1, ll m1, ll x2, ll m2){
  ll g = gcd(m1, m2);
  if ((x2 - x1) % g) return {0, 0}; // no sol
  m1 /= g; m2 /= g;
  11 _, p, q;
  extgcd(m1, m2, _, p, q); // p <= C

ll lcm = m1 * m2 * g;

ll res = ((__int128)p * (x2 - x1) % lcm * m1 % lcm +
       x1) % lcm;
  // be careful with overflow, C^3
  return {(res + lcm) % lcm, lcm}; // (x, m)
}
```

6.11 Sqrt Decomposition [8d7bc0]

```
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){
 int x = ifloor(n, 1);
  r = ifloor(n, x);
// for all i in [l, r], ceil(n / i) = x
for(int 1, r = n; r >= 1; r = 1 - 1){
 int x = iceil(n, r);
  l = iceil(n, x);
```

6.12 Floor Sum

```
• m = \lfloor \frac{an+b}{2} \rfloor
```

• Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

7 **Polynomial**

7.1 FWHT [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)</pre>
    for (int i = 0; i < n; i += L)</pre>
      for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
    ];
void subset_convolution(int *a, int *b, int *c, int L)
  // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
   ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
   f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
  fwt(f[i], n, 1), fwt(g[i], n, 1);
for (int i = 0; i <= L; ++i)</pre>
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i) fwt(h[i], n, -1);</pre>
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
7.2 FFT [13ec2f]
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double
using val_t = complex<double>;
template<int MAXN>
struct FFT {
  const double PI = acos(-1);
  val_t w[MAXN];
  FFT() {
    for (int i = 0; i < MAXN; ++i) {</pre>
      double arg = 2 * PI * i / MAXN;
      w[i] = val_t(cos(arg), sin(arg));
   }
  }
  void bitrev(vector<val_t> &a, int n) //same as NTT
  void trans(vector<val_t> &a, int n, bool inv = false)
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
          val_t tmp = a[j + dl] * (inv ? conj(w[x]) : w
               [x]);
          a[j + dl] = a[j] - tmp;
          a[j] += tmp;
        }
      }
    }
    if (inv) {
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
  //multiplying two polynomials A * B:
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
 //A[i] *= B[i], fft.trans(A, siz, 1);
7.3 NTT [bf683f]
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
// only works when sz(A) + sz(B) - 1 <= MAXN
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
struct NTT {
  11 w[MAXN];
  11 mpow(ll a, ll n);
```

```
11 minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw</pre>
         % P;
  void bitrev(vector<ll> &a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {
  for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
  void operator()(vector<ll> &a, int n, bool inv =
      false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + dl; ++j, x += dx
          ll tmp = a[j + dl] * w[x] % P;
          if ((a[j + d1] = a[j] - tmp) < 0) a[j + d1]
               += P:
          if ((a[j] += tmp) >= P) a[j] -= P;
        }
      }
    if (inv) {
      reverse(a.begin()+1, a.begin()+n);
      11 invn = minv(n);
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
 }
};
7.4 Polynomial Operation [77a8a8]
// Copy from 8BQube
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
    ++i)
#define neg(x) (x ? P - x : 0)
#define V (*this)
template <int MAXN, 11 P, 11 RT> // MAXN = 2^k
struct Poly : vector<ll> {
                                 // coefficients in [0,
  using vector<ll>::vector;
  static inline NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
      .data(), min(p.n(), m), data()); }
  Poly &irev() { return reverse(data(), data() + n()),
      V; }
  Poly &isz(int m) { return resize(m), V; }
  static ll minv(ll x) { return ntt.minv(x); }
  Poly &iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
    return V;
  Poly &imul(ll k) {
    fi(0, n()) V[i] = V[i] * k % P;
    return V;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    assert(m <= MAXN);</pre>
```

Poly X(V, m), Y(rhs, m);

fi(0, m) X[i] = X[i] * Y[i] % P;

return X.isz(n() + rhs.n() - 1);

if (n() == 1) return {minv(V[0])};

while (m < n() * 2) m <<= 1;</pre>

Poly Inv() const $\{ // V[0] != 0, 2*sz \le MAXN \}$

Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);

ntt(X, m), ntt(Y, m);

ntt(X, m, true);

assert(m <= MAXN);</pre>

ntt(Xi, m), ntt(Y, m);

int m = 1;

Poly Y(V, m);

```
fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  ntt(Xi, m, true);
  return Xi.isz(n());
Poly &shift_inplace(const 11 &c) { // 2 * sz <= MAXN
  int n = V.n();
  vector<ll> fc(n), ifc(n);
  fc[0] = ifc[0] = 1;
  for (int i = 1; i < n; i++) {</pre>
    fc[i] = fc[i - 1] * i % P;
    ifc[i] = minv(fc[i]);
  for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
  Poly g(n);
  11 cp = 1;
  for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
       cp = cp * c % P;
  V = V.irev().Mul(g).isz(n).irev();
  for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
      Ρ;
  return V;
Poly shift(const 11 &c) const { return Poly(V).
    shift_inplace(c); }
Poly _Sqrt() const { // Jacobi(V[0], P) = 1
  if (n() == 1) return {QuadraticResidue(V[0], P)};
  Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
      1);
Poly Sqrt() const { // 2 * sz <= MAXN
  Poly a;
  bool has = 0;
  for (int i = 0; i < n(); i++) {</pre>
    if (V[i]) has = 1;
    if (has) a.push_back(V[i]);
  if (!has) return V;
  if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
    return Poly();
  }
  a = a.isz((n() + a.n()) / 2)._Sqrt();
  int sz = a.n();
  a.isz(n());
  rotate(a.begin(), a.begin() + sz, a.end());
  return a;
pair<Poly, Poly> DivMod(const Poly &rhs) const { //
   (rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, V};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs);
 X.irev().isz(m);
  Poly Y(V);
  Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
 X = rhs.Mul(Q), Y = V;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
 Poly ret(n() + 1);
  fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
  return ret;
Poly _tmul(int nn, const Poly &rhs) const {
 Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const vector<ll> &x, const vector<
    Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
```

```
vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
  down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
      _tmul(m, V);
  fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
      1, down[i / 2]);
  vector<11> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return v;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = {neg(x[i]), 1};
  for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
     Mul(up[i * 2 + 1]);
  return up;
vector<ll> Eval(const vector<ll> &x) const { // 1e5,
    1s
  auto up = _tree1(x);
 return _eval(x, up);
static Poly Interpolate(const vector<11> &x, const
    vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _{tree1(x), down(m * 2)};
  vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m - 1; i > 0; --i)
    down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
        down[i * 2 + 1].Mul(up[i * 2]));
 return down[1];
Poly Ln() const { //V[0] == 1, 2*sz <= MAXN
 return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { //V[0] == 0,2*sz <= MAXN
  if (n() == 1) return {1};
  Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln();
  Y[0] = P - 1;
  fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
  return X.Mul(Y).isz(n());
// M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(11 k) const { // 2*sz<=MAXN
  int nz = 0;
  while (nz < n() && !V[nz]) ++nz;</pre>
  if (nz * min(k, (ll)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly{1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const ll c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
      n()).irev();
// sum_j w_j [x^j] f(x^i) for i \in [0, m]
Poly power_projection(Poly wt, int m) { // 4*sz <=
    M\Delta XNI
  assert(n() == wt.n());
  if (!n()) {
    return Poly(m + 1, 0);
  if (V[0] != 0) {
    11 c = V[0];
    V[0] = 0;
    Poly A = V.power_projection(wt, m);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
       factorial
    Poly B(m + 1);
    11 pow = 1;
    fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
        * c % P; // inv. of fac
    A = A.Mul(B).isz(m + 1);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P;
    return A;
  int n = 1;
```

```
while (n < V.n()) n *= 2;</pre>
  isz(n), wt.isz(n).irev();
  int k = 1;
  Poly p(wt, 2 * n), q(V, 2 * n);
  q.imul(P - 1);
  while (n > 1) {
    Poly r(2 * n * k);
    fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
        q[i]));
    Poly pq = p.Mul(r).isz(4 * n * k);
    Poly qq = q.Mul(r).isz(4 * n * k);
    fi(0, 2 * n * k) {
      pq[2 * n * k + i] += p[i];
      qq[2 * n * k + i] += q[i] + r[i];
      pq[2 * n * k + i] %= P;
      qq[2 * n * k + i] \% = P;
    fill(p.begin(), p.end(), 0);
    fill(q.begin(), q.end(), 0);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
   p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];
   q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];
    n /= 2, k *= 2;
  Poly ans(k);
  fi(0, k) ans[i] = p[2 * i];
  return ans.irev().isz(m + 1);
Poly FPSinv() {
  const int n = V.n() - 1;
  if (n == -1) return {};
  assert(V[0] == 0);
  if (n == 0) return V;
  assert(V[1] != 0);
  ll c = V[1], ic = minv(c);
  imul(ic);
  Poly wt(n + 1);
  wt[n] = 1;
  Poly A = V.power_projection(wt, n);
  Poly g(n);
  fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
  g = g.Pow(neg(minv(n)));
  g.insert(g.begin(), 0);
  11 pow = 1:
  fi(0, g.n()) g[i] = g[i] * pow % P, pow = pow * ic
     % P;
  return g;
Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
  return Poly(*this).irev().Mul(rhs).isz(n()).irev();
Poly FPScomp(Poly g) { // solves V(g(x))
  auto rec = [&](auto &rec, int n, int k, Poly Q) ->
      Poly {
    if (n == 1) {
      Poly p(2 * k);
      irev();
      fi(0, k) p[2 * i] = V[i];
      return p;
    Poly R(2 * n * k);
    fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
        Q[i]));
    Poly QQ = Q.Mul(R).isz(4 * n * k);
fi(0, 2 * n * k) {
      QQ[2 * n * k + i] += Q[i] + R[i];
      QQ[2 * n * k + i] %= P;
    Poly nxt_Q(2 * n * k);
for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
      nxt_{Q[n * j + i]} = QQ[(2 * n) * j + (2 * i + 0)
    Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
    Poly pq(4 * n * k);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {</pre>
```

```
pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +
        i];
pq[(2 * n) * j + (2 * i + 1)] %= P;
      Poly p(2 * n * k);
      fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i
          ]) % P;
      pq.pop_back();
      Poly x = pq.TMul(R);
      fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
      return p;
    };
    int sz = 1:
    while(sz < n() || sz < g.n()) sz <<= 1;
    return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
         sz)).isz(sz).irev();
  }
}:
#undef fi
#undef V
#undef neg
using Poly_t = Poly<1 << 19, 998244353, 3>;
```

7.5 Generating Function **Ordinary Generating Function**

- C(x) = A(rx): $c_n = r^n a_n$ 的一般生成函數。
- C(x) = A(x) + B(x): $c_n = a_n + b_n$ 的一般生成函數。
- C(x) = A(x)B(x): $c_n = \sum_{i=0}^n a_i b_{n-i}$ 的一般生成函數。
 $C(x) = A(x)^k$: $c_n = \sum_{i_1+i_2+\ldots+i_k=n} a_{i_1}a_{i_2}\ldots a_{i_k}$ 的一般生成函數。
- C(x) = xA(x)': $c_n = na_n$ 的一般生成函數。
- $C(x) = \frac{A(x)}{1-x}$: $c_n = \sum_{i=0}^n a_i$ 的一般生成函數。
- $C(x) = A(1) + x \frac{A(1) A(x)}{1 x}$: $c_n = \sum_{i=1}^{\infty} a_i$ 的一般生成函數。

常用展開式

- $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$, ${a \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

• 卡特蘭數: $f(x) = \frac{1 - \sqrt{1 - 4x}}{2}$

Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數
- $:\sum_{i_1+i_2+\cdots+i_k=n}^{\sum_{k=0}^{n}\binom{i_1}{n}a_ia_{i_2}\dots a_{i_k}}$ 的指數生成函數 • $\hat{C}(x) = \hat{A}(x)^k$:
- $\hat{C}(x) = \exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 F(0),G(0)=0, $F'(0),G'(0)\neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$
$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

Bostan Mori [41c3bc] 7.6

```
const 11 mod = 998244353;
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = \{P, Q\};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
```

/* center i: radius z[i * 2 + 1] / 2

center i, i + 1: radius z[i * 2 + 2] / 2

```
vector<11> Qn(sz);
                                                                 both aba, abba have radius 2 */
  vector N(2, vector<ll>(sz));
                                                              vector<int> manacher(const string &tmp){ // 0-based
                                                                string s = "%";
  while(k) {
    fill(iter(Qn), 0);
for(int i = 0; i < d + 1; i++){
                                                                int 1 = 0, r = 0;
                                                                for(char c : tmp) s += c, s += '%';
      Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);
                                                                vector<int> z(SZ(s));
                                                                for(int i = 0; i < SZ(s); i++){
  z[i] = r > i ? min(z[2 * l - i], r - i) : 1;
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    }
                                                                  while(i - z[i] >= 0 \&\& i + z[i] < SZ(s)
    ntt(Qn, sz, false);
                                                                         && s[i + z[i]] == s[i - z[i]])
                                                                    ++z[i];
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){
                                                                  if(z[i] + i > r) r = z[i] + i, l = i;
      fill(iter(N[i]), 0);
                                                                }
      copy(iter(M[i]), N[i].begin());
                                                                return z;
                                                             }
      ntt(N[i], sz, false);
      for(int j = 0; j < sz; j++)</pre>
                                                             8.3
                                                                    Lyndon Factorization [7c612b]
        N[i][j] = N[i][j] * Qn[j] % mod;
      ntt(N[i], sz, true);
                                                              // partition s = w[0] + w[1] + ... + w[k-1],
      for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
                                                              // w[0] >= w[1] >= ... >= w[k-1]
        M[i][j >> 1] = N[i][j];
                                                              // each w[i] strictly smaller than all its suffix
                                                              void duval(const string &s, vector<pii> &w) {
                                                                for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
    k >>= 1;
                                                                  for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
  }
                                                                    k = (s[k] < s[j] ? i : k + 1);
  return M[0][0] * ntt.minv(M[1][0]) % mod;
                                                                  // if (i < n / 2 && j >= n / 2) {
                                                                  // for min cyclic shift, call duval(s + s)
11 LinearRecursion(vector<ll> a, vector<ll> c, ll k) {
                                                                  // then here s.substr(i, n / 2) is min cyclic shift
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
                                                                  for (; i <= k; i += j - k)</pre>
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
                                                                    w.pb(pii(i, j - k)); // s.substr(l, len)
      1)));
                                                             }
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
                                                             8.4 Suffix Array [cd67ea]
  auto A = a; A.resize(sz);
                                                              struct SuffixArray {
  auto C = c; C.resize(sz);
                                                                vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]
  ntt(A, sz, false), ntt(C, sz, false);
                                                                    and sa[i-1]
  for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
                                                                                             // sa[0] = s.size()
  ntt(A, sz, true);
                                                                                             // character should be 1-
  A.resize(d);
                                                                                                 hased
                                                                SuffixArray(string& s, int lim=256) { // or
  return BostanMori(A, c, k);
                                                                    basic_string<int>
                                                                  int n = s.size() + 1, k = 0, a, b;
                                                                  vector<int> x(n, 0), y(n), ws(max(n, lim));
     String
8
                                                                  rank.assign(n, 0);
                                                                  for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
      KMP Algorithm [c8b75f]
                                                                  sa = lcp = y, iota(sa.begin(), sa.end(), 0);
                                                                  for (int j = 0, p = 0; p < n; j = max(1, j * 2),
// 0-based
                                                                      lim = p) {
// fail[i] = max k<i s.t. s[0..k] = s[i-k..i]
                                                                    p = j, iota(y.begin(), y.end(), n - j);
vector<int> kmp_build_fail(const string &s){
                                                                    for (int i = 0; i < n; i++)</pre>
  int n = SZ(s);
                                                                      if (sa[i] >= j) y[p++] = sa[i] - j;
  vector<int> fail(n, -1);
                                                                    for (int &i : ws) i = 0;
  int cur = -1;
                                                                    for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
  for(int i = 1; i < n; i++){</pre>
                                                                    for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];</pre>
    while(cur != -1 && s[cur + 1] != s[i])
      cur = fail[cur];
                                                                    swap(x, y), p = 1, x[sa[0]] = 0;
    if(s[cur + 1] == s[i])
                                                                    for(int i = 1; i < n; i++){</pre>
      cur++;
                                                                      a = sa[i - 1], b = sa[i];
    fail[i] = cur;
                                                                      x[b] = (y[a] == y[b] && y[a + j] == y[b + j])?
                                                                            p - 1 : p++;
  return fail;
                                                                  for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
void kmp_match(const string &s, const vector<int> &fail
                                                                  for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
    , const string &t){
                                                                    for (k \&\& k--, j = sa[rank[i] - 1];
  int cur = -1;
                                                                        s[i + k] == s[j + k]; k++);
  int n = SZ(s), m = SZ(t);
                                                                }
  for(int i = 0; i < m; i++){</pre>
                                                             };
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t
        [i]))
                                                              8.5 Suffix Automaton [016373]
      cur = fail[cur];
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
                                                              struct exSAM {
                                                                const int CNUM = 26;
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
                                                                // len: maxlength, link: fail link
  }
                                                                // LenSorted: topo order, cnt: occur
                                                                vector<int> len, link, lenSorted, cnt;
8.2 Manacher Algorithm [caf0f4]
                                                                vector<vector<int>> next;
                                                                int total = 0;
```

int newnode() { return total++;

const int n = s.size();

if (n == 1) return;

```
void init(int n) { // total number of characters
    len.assign(2 * n, 0); link.assign(2 * n, 0);
    lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
next.assign(2 * n, vector<int>(CNUM));
    newnode(), link[0] = -1;
  int insertSAM(int last, int c) {
    // not exSAM: cur = newnode(), p = Last
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len[p] + 1 == len[q]) return link[cur] = q, cur
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[clone][i] = len[next[q][i]] ? next[q][i] :
          0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  }
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue<int> q;
                                                             }
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    vector<int> lc(total);
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
        i]]] = i;
  }
  void solve() {
    for (int i = total - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
8.6 Z-value Algorithm [488d87]
// z[i] = max \ k \ s.t. \ s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z_function(const string &s){
  int n = s.size();
  vector<int> z(n);
  for(int i = 1, l = 0, r = 0; i < n; i++){</pre>
    if(i <= r) z[i] = min(r - i + 1, z[i - 1]);</pre>
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
      z[i]++;
    if(i + z[i] - 1 > r)
      \hat{l} = i, r = i + z[i] - 1;
  return z;
8.7 Main Lorentz [fcfb8f]
struct Rep{ int minl, maxl, len; };
vector<Rep> rep; // 0-base
```

 $// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)$ void main_lorentz(const string &s, int sft = 0) {

```
const int nu = n / 2, nv = n - nu;
  const string u = s.substr(0, nu), v = s.substr(nu),
         ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
             ());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const auto z1 = z_function(ru), z2 = z_function(v +
              z3 = z_function(ru + '#' + rv), z4 =
                  z_function(v);
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
  auto add_rep = [&](bool left, int c, int l, int k1,
      int k2) {
    const int L = max(1, 1 - k2), R = min(1 - left, k1)
    if (L > R) return;
    if (left) rep.emplace_back(Rep({sft + c - R, sft +
         c - L, 1}));
    else rep.emplace_back(Rep({sft + c - R - l + 1, sft
          + c - L - 1 + 1, 1));
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int 1, k1, k2;
    if (cntr < nu) {</pre>
      1 = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
8.8
     AC Automaton [f529e6]
const int SIGMA = 26;
struct AC_Automaton {
  // child: trie, next: automaton
  vector<vector<int>> child, next;
  vector<int> fail, cnt, ord;
  int total = 0;
  int newnode() {
    return total++;
  void init(int len) { // len >= 1 + total len
    child.assign(len, vector<int>(26, -1));
next.assign(len, vector<int>(26, -1));
    fail.assign(len, -1); cnt.assign(len, 0);
    ord.clear();
    newnode();
  int input(string &s) {
    int cur = 0;
    for (char c : s) {
      if (child[cur][c - 'A'] == -1)
    child[cur][c - 'A'] = newnode();
cur = child[cur][c - 'A'];
    return cur; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(0), fail[0] = -1;
    while(!q.empty()) {
      int R = q.front();
      q.pop(); ord.pb(R);
      for (int i = 0; i < SIGMA; i++)</pre>
        if (child[R][i] != -1) {
           int X = next[R][i] = child[R][i], Z = fail[R
           while (Z != -1 && child[Z][i] == -1)
            Z = fail[Z];
           fail[X] = Z != -1 ? child[Z][i] : 0;
           q.push(X);
        else next[R][i] = R ? next[fail[R]][i] : 0;
```

```
}
 }
  void solve() {
    for (int i : ord | views::reverse)
      if (i) cnt[fail[i]] += cnt[i];
};
```

Palindrome Automaton [8a071b] 8.9

```
struct PalindromicTree {
  struct node {
    int nxt[26], fail, len; // num = depth of fail link
    int cnt, num; // cnt = occur, num = #pal_suffix of
        this node
    node(int 1 = 0) : nxt{},fail(0),len(1),cnt(0),num
        (0) \{ \}
 };
  vector<node> st; vector<int> s; int last, n;
  void init() {
    st.clear(); s.clear(); last = 1; n = 0;
    st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
  int getFail(int x) {
    while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x:
  void add(int c) {
    s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
     int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
     st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
  void dpcnt() {
    for(int i = SZ(st) - 1; i >= 0; i--){
     auto nd = st[i];
      st[nd.fail].cnt += nd.cnt;
  int size() { return (int)st.size() - 2; }
};
```

9 Misc

Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int 1 = 0, r = n; bool rv = pred(1, 0);
  while (r - l > 1) {
    int m = (1 + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else 1 = m;
  return pred(1, r % n) ? 1 : r % n;
```

9.2 Matroid

 $M=(E,\mathcal{I})$, where $\mathcal{I}\subseteq 2^E$ is nonempty, is a matroid if:

• If $S \in \mathcal{I}$ and $S' \subsetneq S$, then $S' \in \mathcal{I}$.

• For $S_1, S_2 \in \mathcal{I}$ s.t. $|S_1| < |S_2|$, there exists $e \in S_2 \setminus S_1$ s.t. $S_1 \cup \{e\} \in \mathcal{I}$. Matroid intersection:

Start from $S = \emptyset$. In each iteration, let

• $Y_1 = \{ x \notin S \mid S \cup \{x\} \in \mathcal{I}_1 \}$

• $Y_2 = \{ x \notin S \mid S \cup \{x\} \in \mathcal{I}_2 \}$

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \notin S$, create edges

• $x \to y \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_1.$

• $y \to x \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_2.$

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of ${\cal S}$ will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x)if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.3 Simulate Annealing [ff826c]

```
ld anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<ld> rnd(0, 1);
  const ld dT = 0.001:
  // Argument p
  ld S_cur = calc(p), S_best = S_cur;
  for (ld T = 2000; T > eps; T -= dT) {
    // Modify p to p_prime
    const ld S_prime = calc(p_prime);
    const ld delta_c = S_prime - S_cur;
    ld prob = min((ld)1, exp(-delta_c / T));
    if (rnd(rnd_engine) <= prob)</pre>
      S_cur = S_prime, p = p_prime;
    if (S_prime < S_best) // find min</pre>
      S_best = S_prime, p_best = p_prime;
  return S_best;
```

9.4 Binary Search On Fraction [f6b9ec]

```
struct Q {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p * d, q + b.q * d};
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N, auto &&pred) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step /= 2 : step *=</pre>
        2);)
      if (Q mid = hi.go(lo, len + step);
          mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

9.5 Min Plus Convolution [09b5c3]

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector
     <int> &b) {
  int n = SZ(a), m = SZ(b);
  vector<int> c(n + m - 1, INF);
auto dc = [&](auto Y, int l, int r, int jl, int jr) {
    if (1 > r) return;
    int mid = (1 + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)
  if (int i = mid - j; i >= 0 && i < n)</pre>
         if (best > a[i] + b[j])
           best = a[i] + b[j], from = j;
    Y(Y, l, mid - 1, jl, from), Y(Y, mid + 1, r, from)
         jr);
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
```

10 Notes

10.1 Geometry **Rotation Matrix**

```
/cos θ
              -\sin\theta
\sin \theta
               \cos \theta
```

- rotate 90° : $(x,y) \rightarrow (-y,x)$
- rotate -90° : $(x,y) \rightarrow (y,-x)$

Triangles

Side lengths: a,b,c

Semiperimeter: $p = \frac{a+b+c}{c}$

Area:
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{p}$

Length of median (divides triangle into two equal-area triangles): $m_a =$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}$$

Law of sines:
$$\frac{\sin\alpha}{a}=\frac{\sin\beta}{b}=\frac{\sin\gamma}{c}=\frac{1}{2R}$$
 Law of cosines:
$$a^2=b^2+c^2-2bc\cos\alpha$$

Law of tosines.
$$a = b + c = 2ac \cos \alpha$$

$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \operatorname{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \operatorname{atan2}(u, x) \end{array}$$

Green's Theorem

$$\iint_{D} \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy = \oint_{L^{+}} (Pdx + Qdy)$$

$$\mathsf{Area} = \frac{1}{2} \oint_{I} x \, dy - y \, dx$$

Circular sector:

$$\begin{aligned} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r(r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{aligned}$$

Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1, p_2, p_3 are collinear $\iff p_1^*, p_2^*, p_3^*$ intersect at a point p lies above $l \iff l^*$ lies above p^*
- lower convex hull ↔ upper envelope

Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x})$$

$$\cosh x = \frac{1}{2}(e^x + e^{-x})$$

$$\cos n\pi = (-1)^n$$

$$\begin{split} \sin(\alpha+\beta) &= \sin\alpha\cos\beta + \cos\alpha\sin\beta\\ \cos(\alpha+\beta) &= \cos\alpha\cos\beta - \sin\alpha\sin\beta\\ \sin(2\alpha) &= 2\cos\alpha\sin\alpha\\ \cos(2\alpha) &= \cos^2\alpha - \sin^2\alpha\\ &= 2\cos^2\alpha - 1\\ &= 1 - 2\sin^2\alpha \end{split}$$

$$\begin{split} \tan(\alpha+\beta) &= \frac{\tan\alpha + \tan\beta}{1-\tan\alpha\tan\beta} \\ \sin\alpha + \sin\beta &= 2\sin\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \cos\alpha + \cos\beta &= 2\cos\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \sin\alpha\sin\beta &= \frac{1}{2}(\cos(\alpha-\beta)-\cos(\alpha+\beta)) \\ \sin\alpha\cos\beta &= \frac{1}{2}(\sin(\alpha+\beta)+\sin(\alpha-\beta)) \\ \cos\alpha\sin\beta &= \frac{1}{2}(\sin(\alpha+\beta)-\sin(\alpha-\beta)) \\ \cos\alpha\cos\beta &= \frac{1}{2}(\cos(\alpha-\beta)+\cos(\alpha+\beta)) \\ (V+W)\tan(\alpha-\beta)/2 &= (V-W)\tan(\alpha+\beta)/2 \end{split}$$

where V,W are lengths of sides opposite angles α,β .

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$

where
$$r = \sqrt{a^2 + b^2}, \phi = \text{atan2}(b, a)$$
.

10.3 Calculus

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\begin{split} \frac{d}{dx} \arcsin x &= \frac{1}{\sqrt{1-x^2}} & \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx} \tan x &= 1 + \tan^2 x & \frac{d}{dx} \arctan x = \frac{1}{1+x^2} \\ \int \tan ax &= -\frac{\ln|\cos ax|}{a} & \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2} \\ \int e^{-x^2} &= \frac{\sqrt{\pi}}{2} \mathrm{erf}(x) & \int x e^{ax} &= \frac{e^{ax}}{a^2} (ax-1) \\ \int \sin^2(x) &= \frac{x}{2} - \frac{1}{4} \sin 2x & \int \sin^3 x &= \frac{1}{12} \cos 3x - \frac{3}{4} \cos x \\ \int \cos^2(x) &= \frac{x}{2} + \frac{1}{4} \sin 2x & \int \cos^3 x &= \frac{1}{12} \sin 3x + \frac{3}{4} \sin x \\ \int x \sin x &= \sin x - x \cos x & \int x \cos x &= \cos x + x \sin x \\ \int x e^x &= e^x (x-1) & \int x^2 e^x &= e^x (x^2 - 2x + 2) \end{split}$$

$$\int x^2 \sin x = 2x \sin x - (x^2 - 2) \cos x$$

$$\int x^2 \cos x = 2x \cos x + (x^2 - 2) \sin x$$

$$\int e^x \sin x = \frac{1}{2} e^x (\sin x - \cos x)$$

$$\int e^x \cos x = \frac{1}{2} e^x (\sin x + \cos x)$$

$$\int x e^x \sin x = \frac{1}{2} e^x (x \sin x - x \cos x + \cos x)$$

$$\int x e^x \cos x = \frac{1}{2} e^x (x \sin x + x \cos x - \sin x)$$

10.4 Sum & Series

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

10.5 Misc

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

Let D be a $n \times n$ matrix, where $d_{ij}=x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij}=-d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence d_1, d_2, \ldots, d_n for each $\emph{labeled}$ vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.
- Erdős–Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only

if
$$d_1+\cdots+d_n$$
 is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds

for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers
$$a_1 \geq \cdots \geq a_n$$
 and b_1, \ldots, b_n is bigraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i, k)$

holds for every $1 \le k \le n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$

$$\cdots \geq a_n \text{ is digraphic if and only if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1)$$

$$1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have

 $A = \#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1.$

- · Möbius inversion formula

 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$ $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : arcsin(a/r).
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$ $\cos \theta$)²/3.
 - Area $=2\pi rh=\pi(a^2+h^2)=2\pi r^2(1-\cos\theta).$
- · Lagrange multiplier
 - Optimize $f(x_1,\ldots,x_n)$ when k constraints $g_i(x_1,\ldots,x_n)=0$.
 - Lagrangian function
 - $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k) = f(x_1,\ldots,x_n) \sum_{i=1}^k \lambda_i g_i(x_1,\ldots,x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- · Nearest points of two skew lines
 - Line 1 : $v_1 = p_1 + t_1 d_1$
- Line 2 : $v_2 = p_2 + t_2 d_2$

$$- n = d_1 \times d_2$$

$$- n_1 = d_1 \times n$$

$$- n_2 = d_2 \times n$$

$$- c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1$$

$$egin{align*} oldsymbol{-} & oldsymbol{n} = oldsymbol{d}_1 imes oldsymbol{d}_2 \ oldsymbol{-} & oldsymbol{n}_1 = oldsymbol{d}_1 imes oldsymbol{n}_2 \ oldsymbol{-} & oldsymbol{c}_1 = oldsymbol{p}_1 + rac{(oldsymbol{p}_2 - oldsymbol{p}_1) \cdot oldsymbol{n}_2}{oldsymbol{d}_1 \cdot oldsymbol{n}_2} oldsymbol{d}_1 \ oldsymbol{-} & oldsymbol{c}_2 = oldsymbol{p}_2 + rac{(oldsymbol{p}_2 - oldsymbol{p}_1) \cdot oldsymbol{n}_2}{oldsymbol{d}_2 \cdot oldsymbol{n}_1} oldsymbol{d}_2 \ egin{align*} oldsymbol{P}_2 & oldsymbol{c}_1 & oldsymbol{c}_2 & oldsymbol{c}_2 & oldsymbol{c}_2 & oldsymbol{c}_1 & oldsymbol{c}_2 & oldsymbol$$

• Bernoulli numbers
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^m \binom{m+1}{j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^{n}$$

$$x^{n} = \sum_{i=0}^{n} S(n,i)(x)_{i}$$
 • Pentagonal number theorem

$$x^n = \sum_{i=0}^n S(n,i)(x)_i$$

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
• Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \ge j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$
10.6 Number

$$E(n, 0) = E(n, n - 1) = 1$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{i} {n+1 \choose i} (k+1-i)^{i}$$

· Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039, 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

Number of paritions of n:

Maximum number of divisors:

n | 100 1e3 1e6 1e9 1e12 1e15

|d(i)| 12 32 240 1344 6720 26880 103680

n 1234567 8 10 11 12 13 14 $\binom{2n}{n}$ 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

• Number of ways to partition a set of n labeled elements:

• Fibonacci numbers: $\frac{n}{F_n}$ 1 1 3 3 4 5 31 45 88 1346269 1e9 1e18

