void debug(){cerr << "\n";}
template<class T, class ... U>

void debug(T a, U ... b){cerr << a << " ", debug(b...)</pre> **Contents** 5.11 Minimum Steiner Tree . . . 14 5.12 Count Cycles 15 template < class T> void pary(T 1, T r){ Basic 6 Math 1.1 Default Code while (1 != r) cerr << *1 << " 6.1 Extended Euclidean Algocerr << "\n";</pre> 1.2 .vimrc rithm 15 Fast IO 13 6.2 Floor & Ceil 15 1.4 Random #else 1.5 PBDS Tree 6.3 Legendre 15 #define debug(...) void() 1.6 Pragma 6.4 Simplex 15 #define pary(...) void() 6.5 Simplex Construction . . . 16 #endif 6.6 DiscreteLog 16 2 Data Structure 2.1 Heavy-Light Decomposition 6.7 Miller Rabin & Pollard Rho . 16 template < class A, class B> 2.2 Link Cut Tree 6.8 XOR Basis 16 ostream& operator<<(ostream& o, pair<A,B> p) { return o << '(' << p.ff << ',' << p.ss << 2.3 6.9 Linear Equation 16 << p.ss << ')'; } 2.4 KD Tree 6.10 Chinese Remainder Theorem 17 6.11 Sqrt Decomposition 17 int main(){ 2.7 Dynamic Convex Hull ios_base::sync_with_stdio(false); cin.tie(0); 7 Polynomial 3 Flow & Matching 3.1 Dinic . . 7.1 FWHT 17 3.2 Bounded Flow 1.2 .vimrc [b4816d] 7.2 FFT 17 MCMF 3.3 7.3 NTT 18 Min Cost Circulation 7.4 Polynomial Operation . . . 18 Gomory Hu se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a 7.5 Generating Function 20 3.6 Stoer Wagner Algorithm . . map <F9> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -Ordinary Generating Func-3.7 Bipartite Matching Wextra -Wshadow -O2 -Dzisk -g -fsanitize=address, 3.8 Kuhn Munkres Algorithm . tion 20 6 Exponential Generating undefined<CR> Max Simple Graph Matching 3.9 3.10 Flow Model Function 20 map <F8> :!./%:r<CR> 7.6 Bostan Mori 20 inoremap {<CR> {<CR>}<ESC>ko Geometry ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space 4.1 Geometry Template 8 String :]' \| md5sum \| cut -c-6 Polar Angle Comparator . . 8.1 KMP Algorithm 21 Minkowski Sum inoremap fj <ESC> 8.2 Manacher Algorithm 21 Intersection of Circle and vnoremap fj <ESC> 8.3 Lyndon Factorization 21 " -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG 8.4 Suffix Array 21 Tangent Line of Circles . . . 4.6 8.5 Suffix Automaton 21 1.3 Fast IO [4f6f0e] 4.7 Intersection of Line and 8.6 Z-value Algorithm 22 Convex Polygon 8.7 Main Lorentz 22 char readchar() { 4.8 Intersection of Line and 8.8 AC Automaton 22 const int N = 1 << 20; Circle 8.9 Palindrome Automaton . . 22 4.9 Point in Circle static char buf[N]; 4.10 Point in Convex static char *p = buf , *end = buf; 4.11 Half Plane Intersection if(p == end) { 4.12 Minimum Enclosing Circle . 9.1 Cyclic Ternary Search . . . 23 if((end = buf + fread(buf , 1 , N , stdin)) == buf) 9.2 Matroid 23 return EOF; 9.3 Simulate Annealing 23 p = buf; 4.15 Delaunay Triangulation . . 10 9.4 Binary Search On Fraction . 23 4.16 Voronoi Diagram 10 9.5 Min Plus Convolution . . . 23 4.17 Polygon Union return *p++; 9.6 SMAWK 23 4.18 Tangent Point to Convex Hull 11 } 9.7 Golden Ratio Search 23 4.19 Heart 11 4.20 Rotating Sweep Line 11 const int buf size = 524288; struct Writer { 10.1 Geometry 24 4.23 Calculate Points in Triangle 11 char buf[buf_size]; int size = 0, ret; void flush() { ret = write(1, buf, size); size = 0; } Rotation Matrix 24 Triangles 24 Graph void _flush(int sz) { if (sz + size > buf_size) flush Quadrilaterals 24 5.2 SCC 12 5.3 2-SAT 12 5.4 Dominator Tree 12 Spherical coordinates . . . 24 (); } Green's Theorem 24 void write_char(char c) { _flush(1); buf[size++] = c; Point-Line Duality 24 10.2 Trigonometry 24 void write_int(int x) { 5.6 Fast DMST 13 10.3 Calculus 24 const int len = 20; 10.4 Sum & Series 24 if (x < 0) buf[size++] = '-', x = -x;</pre> 5.9 Number of Maximal Clique 14 5.10 Minimum Mean Cycle . . . 14 10.5 Misc 25 10.6 Number 25 if (x == 0) buf[size + (ptr++)] = '0'; else for (; x; x /= 10) buf[size + (ptr++)] = '0' + x % 10; Basic reverse(buf + size, buf + size + ptr); 1.1 Default Code [2d2229] size += ptr: }; // remember to call flush #include <bits/stdc++.h> using namespace std; 1.4 Random [4cf9ed] #define iter(v) v.begin(),v.end() mt19937 rng(chrono::system clock::now(). #define SZ(v) int(v.size()) time_since_epoch().count()); #define pb emplace_back #define ff first **1.5 PBDS Tree** [9e57e3] #define ss second #include <bits/extc++.h> using 11 = long long; using namespace __gnu_pbds; using pii = pair<int, int>; using Tree = tree<int, null_type, less<>, rb_tree_tag, using pll = pair<ll, ll>; tree_order_statistics_node_update>; // .find_by_order(x) **#ifdef** zisk // .order_of_key(x)

ofstream o; string c = "red";

public: // SVG svg("test.svg", 0, 0, 100, 100)

1.6 Pragma [6006f6]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)

1.7 SVG Writer [7adcc8]

class SVG {
   void p(string_view s) { o << s; }
   void p(string_view s, auto v, auto... vs) {
     auto i = s.find('$');
     o << s.substr(0, i) << v, p(s.substr(i + 1), vs...)</pre>
```

SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f) { p("<svg xmlns='http://www.w3.org/2000/svg' " "viewBox='\$ \$ \$'>\n" "<style>*{stroke-width:0.5%;}</style>\n", x1, -y2, x2 - x1, y2 - y1); } ~SVG() { p("</svg>\n"); } void color(string nc) { c = nc; } void line(auto x1, auto y1, auto x2, auto y2) { p("kine x1='\$' y1='\$' x2='\$' y2='\$' stroke='\$'/>\n", x1, -y1, x2, -y2, c); } void circle(auto x, auto y, auto r) { p("<circle cx='\$' cy='\$' r='\$' stroke='\$' " "fill='none'/>\n", x, -y, r, c); } void text(auto x, auto y, strings, int w = 12) { "fill='none'/>\n", x, -y, r, c); } void text(auto x, auto y, strings, int w = 12) { "fill='none'/>\n", x, -y, r, c); }

p("<text x='\$' y='\$' font-size='\$px'>\$</text>\n",

2 Data Structure

};

x, -y, w, s); }

2.1 Heavy-Light Decomposition [f2dbca]

```
struct HLD{ // 1-based
  int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
  explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1),
  ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
   }
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  void build(){ dfs(1, 1), cut(1, 1); }
int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // Lca
};
```

2.2 Link Cut Tree [cf4f34]

```
// 1-based
// == 43515a ==
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node() : pa{0}, ch{0, 0}, rev{false}, size{1}, v{},
      sum\{\}, rsum\{\}, sv\{\}, sub\{\}, vir\{\} \{\}
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is_rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false:
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
    cur.rev ^= 1;
// == f05d4f ==
  void rotate(int u) {
    int f = cur.pa, g = o[f].pa, l = is_rch(u);
    if (cur.ch[1 ^ 1]) o[cur.ch[1 ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  void splay(int u) {
    vector<int> stk = {u};
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is_root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    up(u);
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
      cur.vir = cur.vir + o[rc].sub - o[last].sub;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0:
    for (access(u); u; u = lc) down(la = u);
    return la;
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
// == a238c2 ==
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); }
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
```

```
return find_root(u) == find_root(v); }
  void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
#undef rc
};
2.3 Treap [2ac37e]
mt19937 rng(880301);
// == fb4359 ==
struct node {
  11 data; int sz;
  node *1, *r;
  node(11 k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
node pool[1000010]; int pool_cnt = 0;
node *newnode(11 k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
  return b->down(), b->l = merge(a, b->l), b->up(), b;
}
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)</pre>
    a = o, split(o->r, a->r, b, k), a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
// a: size k, b: size n - k
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();</pre>
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
}
// == e9f4d8 ==
node *kth(node *o, ll k) { // 1-based
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, 11 key) { // num of key < key</pre>
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, ll k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, ll k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
```

```
tuple<node*, node*, node*> interval(node *&o, int 1,
    int r) { // 1-based
  node *a, *b, *c; // b: [l, r]
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
  return make_tuple(a, b, c);
}
2.4 KD Tree [375ca2]
namespace kdt {
  int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
  point p[maxn];
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
    function < bool (const point &, const point &) > f =
      [dep](const point &a, const point &b) {
        if (dep & 1) return a.x < b.x;</pre>
        else return a.y < b.y;</pre>
      };
    int m = (1 + r) >> 1;
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
    yl[m] = yr[m] = p[m].y;
    lc[m] = build(1, m, dep + 1);
    if (~lc[m]) {
      x1[m] = min(x1[m], x1[1c[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
    return m;
  bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
        q.y < yl[o] - ds || q.y > yr[o] + ds)
      return false;
    return true;
  long long dist(const point &a, const point &b) {
    return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
  void dfs(
      const point &q, long long &d, int o, int dep = 0)
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x ||</pre>
        !(dep & 1) && q.y < p[o].y) {
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    } else {
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    }
  void init(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
    root = build(0, v.size());
  long long nearest(const point &q) {
    long long res = 1e18;
    dfs(q, res, root);
    return res:
} // namespace kdt
2.5 Leftist Tree [e91538]
struct node {
  11 v, data, sz, sum;
node *1, *r;
  node(ll k)
```

}:

```
: v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a\rightarrow r) \rightarrow V(a\rightarrow l)) swap(a\rightarrow r, a\rightarrow l);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
}
2.6 Convex 1D/1D [a449dd]
template < class T>
struct DynamicHull {
```

```
struct seg { int x, l, r; };
  T f; int C; deque<seg> dq; // range: 1~C
  explicit DynamicHull(T _f, int _C): f(_f), C(_C) {} // max t s.t. f(x, t) >= f(y, t), x < y, maintain max
  int intersect(int x, int y) {
    int 1 = 0, r = C + 1;
    while (1 + 1 < r) {
      int mid = (1 + r) / 2;
      if (f(x, mid) >= f(y, mid)) l = mid;
      else r = mid;
    }
    return 1;
  void push_back(int x) {
    for (int i; !dq.empty() &&
         (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
         ();
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back().r + 1 <= C) dq.pb(seg({x, dq.back().r</pre>
          + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
  }
};
```

Dynamic Convex Hull [b45ebc]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
      rhs.a; }
  bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  bool isect(iterator x, iterator y) {
    if (y == end()) \{ x \rightarrow p = kInf; return 0; \}
    if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf
    else x -> p = iceil(y -> b - x -> b, x -> a - y -> a);
    return x->p >= y->p;
  }
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y =
        erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  11 query(ll x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
```

```
Flow & Matching
3
```

3.1 Dinic [801a71]

```
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
    ll to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
    n = _n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
    g[u].pb(edge{v, cap, 0, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
    dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int cur = q.front(); q.pop();
      for (auto &e : g[cur])
        if (dis[e.to] == -1 && e.flow != e.cap) {
          q.push(e.to);
          dis[e.to] = dis[cur] + 1;
        }
      }
    }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
          g[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0:
  11 maxflow(int _s, int _t) {
         _s; t = _t;
    11 \text{ flow} = 0, df;
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
};
3.2
      Bounded Flow [758826]
```

```
struct BoundedFlow : Dinic {
 vector<ll> tot;
 void init(int _n) {
   Dinic::init(_n + 2);
    tot.assign(n, 0);
 void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
```

```
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
                                                             struct MinCostCirculation { // 0-based, O(VE * ElogC)
                                                              struct edge {
  bool feasible() {
                                                                11 from, to, cap, fcap, flow, cost, rev;
    11 \text{ sum} = 0;
    int vs = n - 2, vt = n - 1;
                                                              int n;
    for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
        add_edge(vs, i, 0, tot[i]), sum += tot[i];
      else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
        g[vs].pop_back(), g[i].pop_back();
      else if(tot[i] < 0)</pre>
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
                                                                }:
    add_edge(_t, _s, 0, INF);
    if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    };
                                                                }
                                                              }
3.3 MCMF [671e14]
struct MCMF { // 0-base
  struct Edge {
   11 from, to, cap, flow, cost, rev;
  int n, s, t;
  vector<vector<Edge>> g;
  vector<Edge*> past;
  vector<ll> dis, up, pot;
                                                                  }
  explicit MCMF(int
                     _n): n(_n), g(n), past(n), dis(n),
      up(n), pot(n) {}
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(Edge{a, b, cap, 0, cost, SZ(g[b])});
    g[b].pb(Edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
  bool BellmanFord() {
    vector<bool> inq(n);
    fill(iter(dis), INF);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
                                                                }
        if (!inq[u]) inq[u] = 1, q.push(u);
                                                              }
      }
    };
                                                                n = n:
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
relax(e.to, d2, min(up[u], e.cap - e.flow), &e)
                                                              }
      }
                                                            };
    }
                                                            3.5
    return dis[t] != INF;
  pair<ll, 1l> solve(int _s, int _t, bool neg = true) {
    s = _s, t = _t; 11 flow = 0, cost = 0;
    if (neg) BellmanFord(), pot = dis;
    for (; BellmanFord(); pot = dis) {
      for (int i = 0; i < n; ++i)</pre>
        if (dis[i] != INF) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], g[e.to][e.rev].flow -= up[t];
                                                            }
      }
    return {flow, cost};
};
                                                              int n;
```

vector<edge*> past; vector<vector<edge>> g; vector<ll> dis; void BellmanFord(int s) { vector<int> inq(n); dis.assign(n, INF); queue<int> q; auto relax = [&](int u, ll d, edge *e) { **if** (dis[u] > d) { dis[u] = d, past[u] = e; **if** (!inq[u]) inq[u] = 1, q.push(u); relax(s, 0, 0); while (!q.empty()) { int u = q.front(); q.pop(), inq[u] = 0;for (auto &e : g[u]) if (e.cap > e.flow) relax(e.to, dis[u] + e.cost, &e); void try_edge(edge &cur) { if (cur.cap > cur.flow) return ++cur.cap, void(); BellmanFord(cur.to); if (dis[cur.from] + cur.cost < 0) {</pre> ++cur.flow, --g[cur.to][cur.rev].flow; for (int i = cur.from; past[i]; i = past[i]->from) { auto &e = *past[i]; ++e.flow, --g[e.to][e.rev].flow; ++cur.cap; void solve(int mxlg) { // mxlg >= log(max cap) for (int b = mxlg; b >= 0; --b) { for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) e.cap *= 2, e.flow *= 2; for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) if (e.fcap >> b & 1) try_edge(e); void init(int _n) { past.assign(n, nullptr); g.assign(n, vector<edge>()); void add_edge(ll a, ll b, ll cap, ll cost) { $g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a)}$ == b)}): g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1}); **Gomory Hu** [82d968] void GomoryHu(Dinic &flow) { // 0-based int n = flow.n; vector<int> par(n); for (int i = 1; i < n; ++i) {</pre> flow.reset(); add_edge(i, par[i], flow.maxflow(i, par[i])); for (int j = i + 1; j < n; ++j)</pre> if (par[j] == par[i] && ~flow.dis[j]) par[j] = i;Stoer Wagner Algorithm [a9917b] struct StoerWagner { // 0-based, 0(V^3) vector<int> vis, del; vector<ll> wei;

Min Cost Circulation [47cf18] 3.4

}:

```
vector<vector<ll>> edge;
  void init(int _n) {
   n = _n;
    del.assign(n, 0);
    edge.assign(n, vector<ll>(n));
  void add_edge(int u, int v, ll w) {
   edge[u][v] += w, edge[v][u] += w;
  void search(int &s, int &t) {
   vis.assign(n, 0); wei.assign(n, 0);
    s = t = -1;
    while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
        if (!del[i] && !vis[i] && mx < wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)</pre>
        if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
   }
  ll solve() {
   11 ret = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
      for (int j = 0; j < n; ++j)</pre>
        edge[x][j] = (edge[j][x] += edge[y][j]);
    return ret:
 }
};
3.7
     Bipartite Matching [5bb9be]
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!~my[e] || (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  }
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < n1; ++i)</pre>
     if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
   }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
   n1 = _n1, nr = _nr;
    g.assign(nl, vector<int>());
```

```
Kuhn Munkres Algorithm [683e0a]
```

```
struct KM \{ // O-based, maximum matching, O(V^3) \}
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll> hl, hr, slk;
  vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    // -INF for perfect matching
    w.assign(n, vector<ll>(n, 0));
    pre.assign(n, 0);
    qu.assign(n, 0);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
        0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
               w[x][y])) {
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !s1k[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
         ; hr.assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  3
};
      Max Simple Graph Matching [907d7c]
```

3.9

```
struct Matching { // 0-based, O(V^3)
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> g;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
   }
  void Blossom(int x, int y, int 1) {
    for (; Find(x) != 1; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
```

```
}
  bool Bfs(int r) {
    iota(iter(fa), 0); fill(iter(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : g[x])
        if (s[u] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
             for (int a = u, b = x, last;
    b!= n; a = last, b = pre[a])
               last = match[b], match[b] = a, match[a] =
                    b;
             return true;
          }
           q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
           int 1 = LCA(u, x);
           Blossom(x, u, 1); Blossom(u, x, 1);
    return false;
  Matching(\textbf{int} \_n) : n(\_n), fa(n + 1), s(n + 1), vis(n
  + 1), pre(n + 1, n), match(n + 1, n), g(n) {} void add_edge(int u, int v)
  { g[u].pb(v), g[v].pb(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

3.10 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
- 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
- 1. Consruct super source S and sink T
- 2. For each edge (x,y,c), connect x o y with (cost,cap)=(c,1) if c>0, otherwise connect y o x with (cost,cap)=(-c,1)
- 3. For each edge with c<0, sum these cost as K , then increase d(y) by 1. decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0,d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
- 3. Connect source $s \to v$, $v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} u 2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
- 1. Let $w'(u,v)=w(u,v)-\mu(u)-\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.

- 2. Find the minimum weight matching M with w' . The answer is $\sum \mu(v) + w'(M)$.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\begin{aligned} \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} &\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

4 Geometry

4.1 Geometry Template [86f0f1]

```
using ld = 11;
using pdd = pair<ld, ld>;
#define X first
#define Y second
// Ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){    return v > eps ? 1 : ( v < -eps ? -1
     : 0); }
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)
     ) <= 0; }
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
       p2) || btw(p4, p1, p2))
    return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&</pre>
    ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
     p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
(4))return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
       pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
```

dot(dp, dq));

```
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                                                  fruit advantages
  return q0 + pdd(cross(r - p0, num), dot(r - p0, num))
       / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
                                                               return 1;
4.2 Polar Angle Comparator [808e89]
                                                             }
// -1: a // b (if same), 0/1: a < b
int cmp(pll a, pll b, bool same = true){
#define is_neg(k) (sgn(k.Y) < 0 || (sgn(k.Y) == 0 &&
    sgn(k.X) < 0))
                                                                 int sign1 ){
  int A = is_neg(a), B = is_neg(b);
                                                               vector<Line> ret;
  if(A != B)
    return A < B;</pre>
  if(sgn(cross(a, b)) == 0)
  return same ? abs2(a) < abs2(b) : -1;</pre>
  return sgn(cross(a, b)) > 0;
}
4.3 Minkowski Sum [b3028c]
void reorder_poly(vector<pdd>& pnts){
  int mn = 0;
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn</pre>
        ].Y && pnts[i].X < pnts[mn].X)) mn = i;</pre>
  rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
}
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
 reorder_poly(P); reorder_poly(Q);
  int psz = P.size(), qsz = Q.size();
                                                               return ret;
                                                             }
  P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
  vector<pdd> ans; int i = 0, j = 0;
                                                             4.7
 while(i < psz || j < qsz){
                                                                   [157258]
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1]-P[i], Q[j + 1]-Q[j]));
    if(t >= 0) i++; if(t <= 0) j++;</pre>
  return ans;
                                                               });
}
4.4 Intersection of Circle and Convex Polygon
double _area(pdd pa, pdd pb, double r){
                                                               int n = SZ(C);
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
```

```
double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
   S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double areaPolyCircle(const vector<pdd> poly,const pdd
    &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)*ori(0,
        poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
}
```

4.5 Intersection of Circles [f7a2fe]

```
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
 pdd o1 = a.0, o2 = b.0;
  double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2), d =
      sqrt(d2);
 if(d < max(r1, r2) - min(r1, r2) || d > r1 + r2)
      return 0;
 pdd u = (o1 + o2) * 0.5 + (o1 - o2) * ((r2 * r2 - r1))
      * r1) / (2 * d2));
  double A = sqrt((r1 + r2 + d) * (r1 - r2 + d) * (r1 +
       r2 - d) * (-r1 + r2 + d));
```

```
pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A / (2 * d2)
p1 = u + v, p2 = u - v;
```

4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
        v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sgn(p1.X - p2.X) == 0 and
        sgn(p1.Y - p2.Y) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
```

Intersection of Line and Convex Polygon

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pil a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  if (cmpL(A) < 0 || cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int 1, int r) {
    for (int t = 1; (1 + 1) % n != r; ) {
      int m = ((1 + r + (1 < r? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(t) ? 1 : r) = m;
    return (1 + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
      +1
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned in the same order as the line hits the
      convex */
  return res;
} // convex cut: (r, l]
```

4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
     a, pdd b) {
  pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
  double s = cross(b - a, c - a), h2 = r * r - s * s /
      abs2(b - a);
  if (sgn(h2) < 0) return {};</pre>
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return {p - h, p + h};
```

4.9 Point in Circle [ecf954]

```
// return q's relation with circumcircle of tri(p[0],p
       [1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
       det += __int128(abs2(p[i]) - abs2(q)) * cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

4.10 Point in Convex [f86640]

4.11 Half Plane Intersection [dfb833]

```
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
     .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(iter(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
      return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
          -1)
      continue;
    while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.
        back()))
      dq.pop_back();
    while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
    dq.pb(p);
  while (SZ(dq) >= 3 \&\& !isin(dq[0], dq[SZ(dq) - 2], dq
      .back()))
    dq.pop_back();
  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
    dq.pop_front();
  return vector<Line>(iter(dq));
```

4.12 Minimum Enclosing Circle [5af6d5]

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
   random_shuffle(iter(pts));
   pdd c = pts[0];
   ld r = 0;
   for(int i = 1; i < SZ(pts); i++){
      if(abs(pts[i] - c) <= r) continue;
      c = pts[i]; r = 0;
   for(int j = 0; j < i; j++){
      if(abs(pts[j] - c) <= r) continue;
}</pre>
```

```
9
      c = (pts[i] + pts[j]) / 2;
      r = abs(pts[i] - c);
      for(int k = 0; k < j; k++){
        if(abs(pts[k] - c) > r)
          tie(c, r) = circumcenter(pts[i], pts[j], pts[
    }
  return {c, r};
4.13 3D Point [badbbd]
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x
      (_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
     pi]
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
    .y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
```

return pdd(dot(p, e1), dot(p, e2)); } Point rotate_around(Point p, double angle, Point axis) { double s = sin(angle), c = cos(angle); Point u = axis / abs(axis);

return u * dot(u, p) * (1 - c) + p * c + cross(u, p)

4.14 ConvexHull3D [156311]

e1 = e1 / abs(e1);

e2 = e2 / abs(e2); Point p = u - a;

}

e2 = e2 - e1 * dot(e2, e1);

```
struct convex_hull_3D {
struct Face {
   int a, b, c;
   Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector<Face> res;
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
   int n = SZ(P);
   if (n <= 2) return; // be careful about edge case</pre>
```

struct Delaunay { // 0-base

vector<int> oidx;

```
// ensure first 4 points are not coplanar
                                                              vector<list<Edge>> head; // result udir. graph
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn
                                                              vector<pll> p;
      (abs2(P[0] - p)) != 0; }));
                                                              Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
  swap(P[2], *find_if(iter(P), [&](auto p) { return sgn
                                                                  head(n), p(n) {
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
                                                                iota(iter(oidx), 0);
  swap(P[3], *find_if(iter(P), [&](auto p) { return sgn
                                                                for (int i = 0; i < n; ++i) head[i].clear();</pre>
      (volume(P[0], P[1], P[2], p)) != 0; }));
                                                                sort(iter(oidx), [&](int a, int b)
  vector<vector<int>> flag(n, vector<int>(n));
                                                                    res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
                                                                for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
  for (int i = 3; i < n; ++i) {</pre>
                                                                divide(0, n - 1);
    vector<Face> next;
    for (auto f : res) {
                                                              void addEdge(int u, int v) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
                                                                head[u].push_front(Edge(v));
                                                                head[v].push_front(Edge(u));
      if (d <= 0) next.pb(f);</pre>
                                                                head[u].begin()->twin = head[v].begin();
      int ff = (d > 0) - (d < 0);
                                                                head[v].begin()->twin = head[u].begin();
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
                                                              void divide(int 1, int r) {
                                                                if (1 == r) return;
    for (auto f : res) {
                                                                if (1 + 1 == r) return addEdge(1, 1 + 1);
      auto F = [\&](int x, int y) {
                                                                int mid = (1 + r) \gg 1, nw[2] = \{1, r\};
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
                                                                divide(l, mid), divide(mid + 1, r);
          next.emplace_back(x, y, i);
                                                                auto gao = [&](int t) {
                                                                  pll pt[2] = {p[nw[0]], p[nw[1]]};
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
                                                                  for (auto it : head[nw[t]]) {
                                                                    int v = ori(pt[1], pt[0], p[it.id]);
if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
    }
    res = next;
                                                                        id]) < abs2(pt[1] - pt[0])))
 }
                                                                      return nw[t] = it.id, true;
bool same(Face s, Face t) {
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
                                                                  return false;
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
                                                                while (gao(0) || gao(1));
                                                                addEdge(nw[0], nw[1]); // add tangent
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
                                                                while (true) {
       return 0;
                                                                  pll pt[2] = {p[nw[0]], p[nw[1]]};
                                                                  int ch = -1, sd = 0;
  return 1;
                                                                  for (int t = 0; t < 2; ++t)
int polygon_face_num() {
                                                                    for (auto it : head[nw[t]])
                                                                      if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
 int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
                                                                            -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
    ans += none_of(res.begin(), res.begin() + i, [&](
                                                                           id])))
                                                                        ch = it.id, sd = t;
        Face g) { return same(res[i], g); });
  return ans;
                                                                  if (ch == -1) break; // upper common tangent
                                                                  for (auto it = head[nw[sd]].begin(); it != head[
double get_volume() {
                                                                      nw[sd]].end(); )
                                                                    if (seg_strict_intersect(pt[sd], p[it->id], pt[
  double ans = 0;
                                                                         sd ^ 1], p[ch]))
  for (auto f : res)
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
                                                                      head[it->id].erase(it->twin), head[nw[sd]].
        ]);
                                                                          erase(it++);
                                                                    else ++it;
 return fabs(ans / 6);
                                                                  nw[sd] = ch, addEdge(nw[0], nw[1]);
double get_dis(Point p, Face f) {
 Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
                                                              }
 };
                                                                   Voronoi Diagram [e4f408]
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
                                                            4.16
      .x) * (p3.z - p1.z);
                                                            // all coord. is even, you may want to call
  double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
                                                                halfPlaneInter after then
      .y) * (p3.x - p1.x);
                                                            vector<vector<Line>> vec;
 double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
                                                            void build_voronoi_line(int n, vector<pll> &pts) {
                                                              Delaunay tool(n, pts); // Delaunay
       * a + b * b + c * c);
                                                              vec.clear(), vec.resize(n);
}
                                                              for (int i = 0; i < n; ++i)
  for (auto e : tool.head[i]) {</pre>
};
// n^2 delaunay: facets with negative z normal of
                                                                  int u = tool.oidx[i], v = tool.oidx[e.id];
// convexhull of (x, y, x^2 + y^2), use a pseudo-point
                                                                  pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v]
// (0, 0, inf) to avoid degenerate case
                                                                       · pts[u]);
4.15 Delaunay Triangulation [6a9916]
                                                                  vec[u].pb(Line(m, m + d));
/* Delaunay Triangulation:
  Given a sets of points on 2D plane, find a
   triangulation such that no points will strictly
                                                            4.17 Polygon Union [9fbf66]
   inside circumcircle of any triangle. */
struct Edge {
  int id; // oidx[id]
                                                            ld rat(pll a, pll b) {
                                                              return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
  list<Edge>::iterator twin;
                                                              // all poly. should be ccw
  Edge(int _id = 0):id(_id) {}
                                                            ld polyUnion(vector<vector<pll>>> &poly) {
                                                              1d res = 0;
};
```

for (auto &p : poly)

for (int a = 0; a < SZ(p); ++a) {</pre>

pll A = p[a], B = p[(a + 1) % SZ(p)];

```
vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
          pll C = q[b], D = q[(b + 1) \% SZ(q)];
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            1d sa = cross(D - C, A - C), sb = cross(D -
                 C, B - C);
            segs.pb(sa / (sa - sb), sgn(sc - sd));
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D
                - C)) > 0) {
            segs.pb(rat(C - A, B - A), 1);
            segs.pb(rat(D - A, B - A), -1);
          }
      sort(iter(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
      res += cross(A, B) * sum;
 return res / 2;
4.18 Tangent Point to Convex Hull [523bc1]
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
4.19 Heart [082d19]
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
    (center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
double m = 2. * (x1 * y2 - y1 * x2);
  pdd center;
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
      y1 - y2)) / m;
  center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
       y2 * y2) / m;
  return center + p0;
pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
      s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
\{ return (p1 + p2 + p3) / 3; \}
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
4.20 Rotating Sweep Line [f5f689]
struct Event {
  pll d; int u, v;
  bool operator<(const Event &b) const {</pre>
    int ret = cmp(d, b.d, false);
    return ret == -1 ? false : ret; } // no tie-break
                                                            };
void rotatingSweepLine(const vector<pll> &p) {
  const int n = SZ(p);
  vector<Event> e; e.reserve(n * (n - 1));
```

for (int i = 0; i < n; i++)</pre>

the event occurs

for (int j = 0; j < n; j++) // pos[i] < pos[j] when

```
if (i != j) e.pb(p[j] - p[i], i, j);
  sort(iter(e));
  vector<int> ord(n), pos(n);
  iota(iter(ord), 0);
  \verb|sort(iter(ord), [\&](int i, int j)| { // initial order}|\\
      return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
          X < p[j].X; \});
  for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
  // initialize
  for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
    // do something
    vector<pii> tmp;
    for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
      tmp.pb(pii(e[j].u, e[j].v));
    sort(iter(tmp), [&](pii x, pii y){
        return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
             ], pos[y.ss]); });
    for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
      tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
        make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
             pos[x]);
4.21 Vector In Poly [c6d0fa]
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
     a-c
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
  return ori(a, b, p) >= strict && ori(a, p, c) >=
// whether vector{cur, p} in counter-clockwise order
    prv, cur, nxt
bool inside(pll prv, pll cur, pll nxt, pll p, int
    strict) {
  if (ori(cur, nxt, prv) >= 0)
    return btwangle(cur, nxt, prv, p, strict);
  return !btwangle(cur, prv, nxt, p, !strict);
4.22 Convex Hull DP [6dc001]
sort(iter(pts), [&](pll x, pll y) {
    return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;</pre>
    }):
auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
    ]; };
vector<pii> trans;
for (int j = 0; j < n; j++)</pre>
    for (int k = 0; k < n; k++)
        if (j != k) trans.pb(pii(j, k));
sort(iter(trans), [\&](pii x, pii y) \rightarrow bool{
    int tmp = cmp(getvec(x), getvec(y), false);
    if (tmp != -1) return tmp;
    pll v = getvec(x);
    return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
});
// DP for convex hull vertices (no points on edges)
auto solve = [\&](int bottom) \{ // O(n^3) \}
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether bottom -> j is legal
    // init trans -> j
  for (auto [i, j] : trans) {
    if (i <= bottom || j <= bottom ||</pre>
            ori(pts[bottom], pts[i], pts[j]) <= 0)</pre>
                 continue;
    // check whether i -> j is legal
    // normal trans i -> j
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether j -> bottom is legal
    // end trans j ->
  }
for(int i = 0; i < n; i++) solve(i);</pre>
      Calculate Points in Triangle [bf746f]
// all points are distinct
```

// cnt[i][j] = # of point k s.t. strictly above ij, and

i < k < j

```
// cnt2[i][j] = # of points k s.t. strictly in ij
// preprocess space: O(n^2), time: O(n^3), query time:
    0(1)
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
  for (int j = 0; j < n; j++){</pre>
    if (pts[i] >= pts[j]) continue;
    for (int k = 0; k < n; k++) {
      if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
        int tmp = ori(pts[i], pts[j], pts[k]);
if (tmp > 0) cnt[i][j]++; // only for i < j</pre>
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
      }
    }
  }
auto calc_tri = [&](array<int, 3> arr) { // strictly
  sort(iter(arr), [&](int x, int y){ return pts[x] <</pre>
      pts[y]; });
  auto [x, y, z] = arr;
  int tmp = ori(pts[x], pts[y], pts[z]);
  if (tmp == 0) return 0;
  else if (tmp < 0)</pre>
    return cnt[x][z] - cnt[x][y] - cnt[y][z] - cnt2[x][
        y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
      x][z];
};
```

5 Graph

5.1 BCC [d04ebe]

```
struct BCC{ // 0-based, allow multi edges but not allow
     Loops
  int n, m, cnt = 0;
  // n:|V|, m:|E|, cnt:#bcc
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
 vector<vector<int>> bcc_v, bcc_e;
vector<vector<pii>> g; // original graph
  vector<pii> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
    n(_n), m(SZ(_edges)), g(_n), edges(_edges){
      for(int i = 0; i < m; i++){</pre>
        auto [u, v] = edges[i];
        g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
    }
  void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
  // modify these if you need more information
void add_v(int v){ bcc_v.back().pb(v); }
  void add_e(int e){ bcc_e.back().pb(e); }
  void build(){
    vector\langle int \rangle in(n, -1), low(n, -1), stk;
    vector<vector<int>> up(n);
    int ts = 0;
    auto _dfs = [&](auto dfs, int now, int par, int pe)
          -> void{
      if(pe != -1) up[now].pb(pe);
      in[now] = low[now] = ts++;
      stk.pb(now):
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
        if(in[v] != -1){
           if(in[v] < in[now]) up[now].pb(e);</pre>
          low[now] = min(low[now], in[v]);
          continue;
        dfs(dfs, v, now, e);
        low[now] = min(low[now], low[v]);
      if((now != par && low[now] >= in[par]) || (now ==
            par && SZ(g[now]) == 0)){
        make_bcc();
        for(int v = stk.back();; v = stk.back()){
          stk.pop_back(), add_v(v);
           for(int e : up[v]) add_e(e);
          if(v == now) break;
        if(now != par) add_v(par);
    };
```

```
12
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) _dfs(_dfs, i, i, -1);
  }
};
5.2 SCC [2c9a01]
struct SCC{ // 0-based, output reversed topo order
  int n, cnt = 0;
  vector<vector<int>> g;
  vector<int> sccid;
  explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
void add_edge(int u, int v){
    g[u].pb(v);
  void build(){
    vector<int> in(n, -1), low(n), stk;
    vector<bool> instk(n);
    int ts = 0;
    auto dfs1 = [&](auto dfs, int now) -> void{
      stk.pb(now); instk[now] = true;
      in[now] = low[now] = ts++;
      for(int i : g[now]){
        if(in[i] == -1)
          dfs(dfs, i), low[now] = min(low[now], low[i])
        else if(instk[i] && in[i] < in[now])</pre>
          low[now] = min(low[now], in[i]);
      if(low[now] == in[now]){
        for(; stk.back() != now; stk.pop_back())
          sccid[stk.back()] = cnt, instk[stk.back()] =
               false;
        sccid[now] = cnt++, instk[now] = false, stk.
             pop_back();
      }
    };
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) dfs1(dfs1, i);
  }
};
5.3 2-SAT [0686a5]
struct SAT { // 0-based
  int n;
  vector<bool> istrue;
  SCC scc;
  SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
  int neg(int a) {
    return a >= n ? a - n : a + n;
  void add_clause(int a, int b) {
    scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
  bool solve() {
    scc.build();
    for (int i = 0; i < n; ++i) {</pre>
      if (scc.sccid[i] == scc.sccid[i + n]) return
           false;
      istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>
      istrue[i + n] = !istrue[i];
    return true;
  }
};
5.4 Dominator Tree [2da9bb]
struct Dominator {
  int n:
  vector<vector<int>> g, r, rdom; int tk;
  vector<int> dfn, rev, fa, sdom, dom, val, rp;
  Dominator(\textbf{int} \_n) : n(\_n), g(n), r(n), rdom(n), tk(0)
    dfn = rev = fa = sdom = dom =
      val = rp = vector<int>(n, -1); }
  void add_edge(int x, int y) { g[x].push_back(y); }
  void dfs(int x) {
    rev[dfn[x] = tk] = x;
    fa[tk] = sdom[tk] = val[tk] = tk; tk++;
    for (int u : g[x]) {
```

if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];

```
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      r[dfn[u]].push_back(dfn[x]);
    }
  }
  void merge(int x, int y) { fa[x] = y; }
  int find(int x, int c = 0) {
    if (fa[x] == x) return c ? -1 : x;
    if (int p = find(fa[x], 1); p != -1) {
      if (sdom[val[x]] > sdom[val[fa[x]]])
        val[x] = val[fa[x]];
      fa[x] = p;
      return c ? p : val[x];
    } else return c ? fa[x] : val[x];
  }
  vector<int> build(int s) {
    // return the father of each node in dominator tree
    dfs(s); // p[i] = -2 if i is unreachable, par[s] =
    for (int i = tk - 1; i >= 0; --i) {
      for (int u : r[i])
        sdom[i] = min(sdom[i], sdom[find(u)]);
      if (i) rdom[sdom[i]].push_back(i);
      for (int u : rdom[i]) {
        int p = find(u);
        dom[u] = (sdom[p] == i ? i : p);
      if (i) merge(i, rp[i]);
    }
    vector < int > p(n, -2); p[s] = -1;
    for (int i = 1; i < tk; ++i)</pre>
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < tk; ++i)</pre>
      p[rev[i]] = rev[dom[i]];
    return p;
 }
};
    Virtual Tree [6abeb5]
5.5
vector<int> vG[N]:
int top, st[N];
int vrt = -1;
void insert(int u) {
  if (top == -1) return st[++top] = vrt = u, void();
  int p = LCA(st[top], u);
    if(dep[vrt] > dep[p]) vrt = p;
  if (p == st[top]) return st[++top] = u, void();
  while (top \Rightarrow 1 && dep[st[top - 1]] \Rightarrow dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
 top = -1;
  sort(iter(v),
      [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do somethina
  reset(vrt);
5.6 Fast DMST [7b274d]
struct E { int s, t; ll w; }; // 0-base
struct PQ {
  struct P {
    11 v; int i;
    bool operator>(const P &b) const { return v > b.v;
  };
  priority_queue<P, vector<P>, greater<>> pq; ll tag;
      // min heap
  void push(P p) { p.v -= tag; pq.emplace(p); }
  P top() { P p = pq.top(); p.v += tag; return p; }
  void join(PQ &b) {
    if (pq.size() < b.pq.size())</pre>
      swap(pq, b.pq), swap(tag, b.tag);
```

while (!b.pq.empty()) push(b.top()), b.pq.pop();

```
}; // O(E log^2 V), use leftist tree for O(E log V)
vector<int> dmst(const vector<E> &e, int n, int root) {
    vector<PQ> h(n * 2);
    for (int i = 0; i < int(e.size()); ++i)</pre>
        h[e[i].t].push({e[i].w, i});
    vector<int> a(n * 2); iota(iter(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
    auto o = [\&](auto Y, int X) \rightarrow int {
        return x==a[x] ? x : a[x] = Y(Y, a[x]); };
    auto S = [&](int i) { return o(o, e[i].s); };
    int pc = v[root] = n;
    for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
        for (int p = i; v[p]<0 \mid \mid v[p]==i; p = S(r[p])) {
            if (v[p] == i)
                for (int q = pc++; p != q; p = S(r[p])) {
                   h[p].tag -= h[p].top().v; h[q].join(h[p]);
                   pa[p] = a[p] = q;
            while (S(h[p].top().i) == p) h[p].pq.pop();
            v[p] = i; r[p] = h[p].top().i;
        }
    vector<int> ans;
    for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
        for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
                1)
            v[f] = n;
        ans.push_back(r[i]);
    return ans; // default minimize, returns edgeid array
5.7 Vizing [58a6ca]
// find D+1 edge coloring of a graph with max deg D, O(
struct Vizing { // returns maxdeg+1 edge coloring in
        adjacent matrix G
    int n; // 1-based for vertices and colors, simple
            araph
    vector<vector<int>> C, G;
    vector<int> X, vst;
    Vizing(int _n): n(_n),
    C(n + 1, vector < int > (n + 2)), G(n + 1, vector < int > (n + 1, 
            + 1)),
    X(n + 1, 1), vst(n + 1) {}
    void solve(vector<pii> &E) {
        auto update = [&](int u)
        { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
        auto color = [&](int u, int v, int c) {
            int p = G[u][v];
            G[u][v] = G[v][u] = c;
            C[u][c] = v, C[v][c] = u;
            C[u][p] = C[v][p] = 0;
            if (p) X[u] = X[v] = p;
            else update(u), update(v);
            return p;
        auto flip = [&](int u, int c1, int c2) {
            int p = C[u][c1];
            swap(C[u][c1], C[u][c2]);
            if (p) G[u][p] = G[p][u] = c2;
            if (!C[u][c1]) X[u] = c1;
            if (!C[u][c2]) X[u] = c2;
            return p;
        for (int t = 0; t < SZ(E); ++t) {</pre>
            int u = E[t].ff, v0 = E[t].ss, v = v0, c0 = X[u],
                      c = c0, d;
            vector<pii> L;
            fill(iter(vst), 0);
            while (!G[u][v0]) {
               L.emplace_back(v, d = X[v]);
                if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
                        --a) c = color(u, L[a].ff, c);
                else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
                          0; --a) color(u, L[a].ff, L[a].ss);
                else if (vst[d]) break;
                else vst[d] = 1, v = C[u][d];
            if (!G[u][v0]) {
                for (; v; v = flip(v, c, d), swap(c, d));
                if (int a; C[u][c0]) {
```

```
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          for (a = SZ(L) - 2; a >= 0 && L[a].ss != c;
               --a):
          for (; a >= 0; --a) color(u, L[a].ff, L[a].ss
        else --t;
     }
   }
 }
};
     Maximum Clique [1ad4b2]
5.8
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(iter(r), [\&](int x, int y) \{ return d[x] > d \}
          [y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
     int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < 1ft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
          ]._Find_next(p))
```

5.9 Number of Maximal Clique [11fa26]

pre_dfs(r, 0, bitset<N>(string(n, '1')));

r[tp] = p, c[tp] = k, ++tp;

void dfs(vector<int> &r, vector<int> &c, int 1,

for (int i : r) if (G[p][i]) nr.pb(i);

if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);

else if (q > ans) ans = q, copy_n(cur, q, sol);

dfs(r, c, l + 1, mask);

bitset<N> mask) {
while (!r.empty()) {

int p = r.back();

c.pop_back(), --q;

cur[q++] = p; vector<int> nr;

vector<int> r(n);

}

}

};

int solve() {

return ans;

r.pop_back(), mask[p] = 0;
if (q + c.back() <= ans) return;</pre>

ans = q = 0, iota(iter(r), 0);

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
        for (int j = 1; j <= n; ++j) g[i][j] = 0;
  }
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  }
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
```

```
int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
int tsn = 0, tnn = 0;
       copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
           some[d + 1][tsn++] = some[d][j];
       for (int j = 0; j < nn; ++j)</pre>
         if (g[v][none[d][j]])
           none[d + 1][tnn++] = none[d][j];
       dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S:
  }
};
```

5.10 Minimum Mean Cycle [3e5d2b]

```
11 road[N][N]; // input here
struct MinimumMeanCycle {
  11 dp[N + 5][N], n;
  pll solve() {
    11 a = -1, b = -1, L = n + 1;
    for (int i = 2; i <= L; ++i)</pre>
      for (int k = 0; k < n; ++k)
        for (int j = 0; j < n; ++j)
          dp[i][j] =
            min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
      if (dp[L][i] >= INF) continue;
      11 ta = 0, tb = 1;
      for (int j = 1; j < n; ++j)</pre>
        if (dp[j][i] < INF &&</pre>
          ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
          ta = dp[L][i] - dp[j][i], tb = L - j;
      if (ta == 0) continue;
      if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      ll g = \_gcd(a, b);
      return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
};
```

5.11 Minimum Steiner Tree [21acea]

```
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcost[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
    }
  }
  void add_edge(int ui, int vi, int wi) {
    dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
  for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
           dst[i][j] =
             min(dst[i][j], dst[i][k] + dst[k][j]);
```

for (; m > 1;) {

if (a == 0) **return** 0;

const int r = __builtin_ctz(a);

if ((r & 1) & & ((m + 2) & 4)) s = -s;

a %= m;

```
int solve(const vector<int> &ter) {
    shortest_path();
                                                                 if (a \& m \& 2) s = -s;
    int t = SZ(ter);
                                                                 swap(a, m);
    for (int i = 0; i < (1 << t); ++i)
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
                                                               return s;
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
                                                             }
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
                                                             // 0: a == 0
        int who = __lg(msk);
for (int i = 0; i < n; ++i)</pre>
                                                             // -1: a isn't a quad res of p
                                                             // else: return X with X^2 \% p == a
                                                             // doesn't work with long long
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
                                                             int QuadraticResidue(int a, int p) {
                                                               if (p == 2) return a & 1;
      for (int i = 0; i < n; ++i)</pre>
                                                               if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
        for (int submsk = (msk - 1) & msk; submsk;
                                                               int b, d;
              submsk = (submsk - 1) \& msk)
                                                               for (; ; ) {
                                                                 b = rand() % p;
d = (1LL * b * b + p - a) % p;
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
                                                                 if (Jacobi(d, p) == -1) break;
              vcost[i]);
      for (int i = 0; i < n; ++i) {</pre>
        tdst[i] = INF;
                                                               int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
        for (int j = 0; j < n; ++j)
                                                               for (int e = (1LL + p) >> 1; e; e >>= 1) {
          tdst[i] =
                                                                 if (e & 1) {
                                                                   tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
            min(tdst[i], dp[msk][j] + dst[j][i]);
                                                                         p)) % p;
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
                                                                   g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
                                                                   g0 = tmp;
    int ans = INF;
                                                                 tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
    for (int i = 0; i < n; ++i)</pre>
                                                                 )) % p;
f1 = (2LL * f0 * f1) % p;
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans:
                                                                 f0 = tmp;
  }
};
                                                               return g0;
5.12 Count Cycles [c7e8f2]
                                                             }
                                                             6.4 Simplex [aa7741]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
                                                             // maximize c^T x
for (int x : ord) { // c3
                                                             // subject to Ax <= b, x >= 0
  for (int y : D[x]) vis[y] = 1;
                                                             // and stores the solution;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
                                                             typedef long double T; // Long double, Rational, double
  for (int y : D[x]) vis[y] = 0;
                                                                   + mod<P>...
                                                             typedef vector<T> vd;
for (int x : ord) { // c4
                                                             typedef vector<vd> vvd;
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) c4 += vis[z]++;
                                                             const T eps = 1e-9, inf = 1/.0;
  for (int y : D[x]) for (int z : adj[y])
                                                             #define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
    if (rk[z] > rk[x]) --vis[z];
                                                                 s])) s=j
} // both are O(M*sqrt(M))
                                                             #define rep(i, l, n) for(int i = l; i < n; i++)
6 Math
                                                             struct LPSolver {
      Extended Euclidean Algorithm [c51ae9]
                                                               int m, n;
                                                               vector<int> N, B;
// ax+ny = 1, ax+ny == ax == 1 \ (mod n)
                                                               vvd D:
void extgcd(ll x, ll y, ll &g, ll &a, ll &b) {
  if (y == 0) g = x, a = 1, b = 0;
                                                               LPSolver(const vvd& A, const vd& b, const vd& c) :
  else extgcd(y, x \% y, g, b, a), b -= (x / y) * a;
                                                                 m(SZ(b)), n(SZ(c)), N(n+1), B(m), D(m+2, vd(n+2)) {
                                                                   rep(i,0,m) \ rep(j,0,n) \ D[i][j] = A[i][j];
                                                                    rep(i,0,m) \{ B[i] = n+i; D[i][n] = -1; D[i][n+1] \}
6.2 Floor & Ceil [134881]
                                                                        = b[i];
                                                                   rep(j,0,n) { N[j] = j; D[m][j] = -c[j]; }
ll ifloor(ll a, ll b){
                                                                   N[n] = -1; D[m+1][n] = 1;
  return a / b - (a % b && (a < 0) ^ (b < 0));
ll iceil(ll a, ll b){
                                                               void pivot(int r, int s) {
  return a / b + (a % b && (a < 0) ^ (b > 0));
                                                                 T *a = D[r].data(), inv = 1 / a[s];
                                                                  rep(i,0,m+2) if (i != r && abs(D[i][s]) > eps) {
                                                                   T *b = D[i].data(), inv2 = b[s] * inv;
rep(j,0,n+2) b[j] -= a[j] * inv2;
b[s] = a[s] * inv2;
6.3 Legendre [4e4b23]
// the Jacobi symbol is a generalization of the
    Legendre symbol,
// such that the bottom doesn't need to be prime.
                                                                 rep(j,0,n+2) if (j != s) D[r][j] *= inv;
// (n|p) -> same as Legendre
                                                                 rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
// (n|ab) = (n|a)(n|b)
                                                                 D[r][s] = inv;
// work with long long
                                                                 swap(B[r], N[s]);
int Jacobi(int a, int m) {
  int s = 1;
```

bool simplex(int phase) {

for (;;) {

int s = -1;

int x = m + phase - 1;

rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);

bool Miller_Rabin(ll a, ll n) {

if ((a = a % n) == 0) return 1; if (n % 2 == 0) return n == 2;

11 tmp = (n - 1) / ((n - 1) & (1 - n));

```
if (D[x][s] >= -eps) return true;
       int r = -1;
                                                                        if (tmp & 1) x = mul(x, a, n);
       rep(i,0,m) {
         if (D[i][s] <= eps) continue;
if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])</pre>
                                                                      if (x == 1 || x == n - 1) return 1;
                                                                      while (--t)
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
                                                                        if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                      return 0;
       if (r == -1) return false;
       pivot(r, s);
                                                                   bool prime(ll n){
                                                                      vector<ll> tmp = {2, 325, 9375, 28178, 450775,
    }
                                                                          9780504, 1795265022};
                                                                      for(ll i : tmp)
  T solve(vd &x) {
                                                                        if(!Miller_Rabin(i, n)) return false;
     int r = 0;
                                                                      return true;
     rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
     if (D[r][n+1] < -eps) {</pre>
                                                                   map<11, int> cnt;
       pivot(r, n);
                                                                   void PollardRho(ll n) {
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                     if (n == 1) return;
                                                                      if (prime(n)) return ++cnt[n], void();
            inf;
       rep(i,0,m) if (B[i] == -1) {
                                                                      if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
                                                                          void();
         rep(j,1,n+1) ltj(D[i]);
                                                                      11 \times 2, y = 2, d = 1, p = 1;
                                                                   #define f(x, n, p) ((mul(x, x, n) + p) % n)
         pivot(i, s);
                                                                      while (true) {
       }
                                                                        if (d != n && d != 1) {
    bool ok = simplex(1); x = vd(n);
                                                                          PollardRho(n / d);
    rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
                                                                          PollardRho(d);
     return ok ? D[m][n+1] : inf;
                                                                          return;
  }
                                                                        if (d == n) ++p;
};
                                                                        x = f(x, n, p), y = f(f(y, n, p), n, p);
       Simplex Construction
                                                                        d = gcd(abs(x - y), n);
Standard form: maximize \sum_{1 < i < n} c_i x_i such that \sum_{1 < i < n} A_{ji} x_i \le b_j for
all 1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
1. In case of minimization, let c_i^\prime = -c_i
                                                                        XOR Basis [006505]
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j \rightarrow \mathsf{add} \subseteq \mathsf{and} \supseteq.
                                                                   const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                   struct Basis{
6.6 DiscreteLog [da27bf]
                                                                      int total = 0, rank = 0;
                                                                      vector<ll> b;
int DiscreteLog(int s, int x, int y, int m) {
                                                                      Basis(): b(digit) {}
  constexpr int kStep = 32000;
                                                                      bool add(l1 v){ // Gauss Jordan Elimination
  unordered_map<int, int> p;
                                                                        total++;
  int b = 1;
                                                                        for(int i = digit - 1; i >= 0; i--){
  for (int i = 0; i < kStep; ++i) {</pre>
                                                                          if(!(1LL << i & v)) continue;</pre>
                                                                          if(b[i] != 0){
    p[y] = i;
    y = 1LL * y * x % m;
                                                                            v ^= b[i];
    b = 1LL * b * x % m;
                                                                            continue;
  for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
                                                                          for(int j = 0; j < i; j++)
  if(1LL << j & v) v ^= b[j];</pre>
                                                                          for(int j = i + 1; j < digit; j++)</pre>
     if (p.find(s) != p.end()) return i + kStep - p[s];
  }
                                                                            if(1LL << i & b[j]) b[j] ^= v;</pre>
  return -1;
                                                                          b[i] = v;
                                                                          rank++;
int DiscreteLog(int x, int y, int m) {
                                                                          return true;
  if (m == 1) return 0;
  int s = 1;
                                                                        return false:
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
s = 1LL * s * x % m;
                                                                      11 \text{ getmax}(11 \text{ x} = 0)
                                                                        for(ll i : b) x = max(x, x ^ i);
                                                                        return x;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
                                                                      ll getmin(ll x = 0){
  if (fpow(x, p, m) != y) return -1;
                                                                        for(ll i : b) x = min(x, x ^ i);
  return p; //returns: x^p = y (mod m)
                                                                        return x:
                                                                      bool can(ll x){
6.7 Miller Rabin & Pollard Rho [d3ecd2]
                                                                        return getmin(x) == 0;
// n < 4,759,123,141 3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
                                                                      11 kth(11 k){ // kth smallest, 0-indexed
                                                                        vector<ll> tmp;
// n < 3,474,749,660,383 6 : primes <= 13
                                                                        for(ll i : b) if(i) tmp.pb(i);
// n < 2^64
                                                                        11 \text{ ans} = 0:
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                        for(int i = 0; i < SZ(tmp); i++)</pre>
11 mul(ll a, ll b, ll n){
                                                                          if(1LL << i & k) ans ^= tmp[i];</pre>
  return (__int128)a * b % n;
```

Linear Equation [056191] 6.9

return ans;

};

```
vector<int> RREF(vector<vector<11>> &mat) { // 9cd26b
 int N = SZ(mat), M = SZ(mat[0]);
  int rk = 0;
  vector<int> cols;
  for (int i = 0; i < M; i++) {</pre>
    int cnt = -1;
    for (int j = N - 1; j >= rk; j--)
     if(mat[j][i] != 0) cnt = j;
    if (cnt == -1) continue;
    swap(mat[rk], mat[cnt]);
    ll lead = mat[rk][i];
    for (int j = 0; j < M; j++) mat[rk][j] = mat[rk][j]</pre>
         * inv(lead) % MOD;
    for (int j = 0; j < N; j++) {</pre>
      if (j == rk) continue;
      11 tmp = mat[j][i];
      for (int k = 0; k < M; k++)
        mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % MOD
             + MOD) % MOD;
    cols.pb(i);
    rk++;
  }
  return cols;
struct LinearEquation { // 2702e2
 bool ok;
  vector<11> par; //particular solution (Ax = b)
  vector<vector<ll>>> homo; //homogenous (Ax = 0)
  vector<vector<ll>> rref;
 //first M columns are matrix A
  //last column of eq is vector b
  void solve(const vector<vector<ll>>> &eq) {
    int M = SZ(eq[0]) - 1;
    rref = eq;
    auto piv = RREF(rref);
    int rk = piv.size();
    if(piv.size() && piv.back() == M)
     return ok = 0, void();
    ok = 1;
    par.resize(M);
    vector<bool> ispiv(M);
    for (int i = 0;i < rk;i++) {</pre>
      par[piv[i]] = rref[i][M];
      ispiv[piv[i]] = 1;
    for (int i = 0; i < M; i++) {</pre>
      if (ispiv[i]) continue;
      vector<ll> h(M);
      h[i] = 1;
      for (int j = 0; j < rk; j++)</pre>
        h[piv[j]] = rref[j][i] ? MOD - rref[j][i] : 0;
      homo.pb(h);
    }
 }
};
6.10 Chinese Remainder Theorem [6ef4a3]
pll solve_crt(ll x1, ll m1, ll x2, ll m2){
 11 g = gcd(m1, m2);
  if ((x2 - x1) % g) return {0, 0}; // no sol
 m1 /= g; m2 /= g;
  11 _, p, q;
 extgcd(m1, m2, _, p, q); // p <= C
ll lcm = m1 * m2 * g;
 ll res = ((__int128)p * (x2 - x1) % lcm * m1 % lcm +
      x1) % lcm;
  // be careful with overflow, C^3
 return {(res + lcm) % lcm, lcm}; // (x, m)
}
6.11 Sqrt Decomposition [8d7bc0]
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){
 int x = ifloor(n, 1);
 r = ifloor(n, x);
}
^{\prime}// for all i in [l, r], ceil(n / i) = x
for(int 1, r = n; r >= 1; r = 1 - 1){
```

int x = iceil(n, r);
l = iceil(n, x);

6.12 Floor Sum

```
• m = \lfloor \frac{an+b}{c} \rfloor
```

• Time complexity: $O(\log n)$

$$f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases}$$

$$g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases}$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \end{cases}$$

$$= \begin{cases} l - \frac{a}{c} \rfloor \cdot (n(n+1)m - f(c, c - b - 1, a, m - 1) \\ -h(c, c - b - 1, a, m - 1)), & \text{otherwise} \end{cases}$$

$$h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^{2}$$

$$= \begin{cases} \lfloor \frac{a}{c} \rfloor^{2} \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^{2} \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \end{cases}$$

$$= \begin{cases} l - \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases}$$

$$nm(m+1) - 2g(c, c - b - 1, a, m - 1)$$

$$-2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases}$$

7 Polynomial

7.1 FWHT [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)</pre>
    for (int i = 0; i < n; i += L)</pre>
       for (int j = i; j < i + (L >> 1); ++j)
         a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
     ];
void subset_convolution(int *a, int *b, int *c, int L)
    {
  // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
  for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
         h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i) fwt(h[i], n,</pre>
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
```

7.2 FFT [13ec2f]

```
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double

using val_t = complex<double>;
template<int MAXN>
struct FFT {
    const double PI = acos(-1);
    val_t w[MAXN];
    FFT() {
       for (int i = 0; i < MAXN; ++i) {</pre>
```

++i)

```
double arg = 2 * PI * i / MAXN;
                                                               #define neg(x) (x ? P - x : 0)
      w[i] = val_t(cos(arg), sin(arg));
                                                               #define V (*this)
                                                               template <int MAXN, ll P, ll RT> // MAXN = 2^k
    }
                                                               struct Poly : vector<ll> { // coefficients in [0, P)
  void bitrev(vector<val_t> &a, int n) //same as NTT
                                                                 using vector<ll>::vector;
  void trans(vector<val_t> &a, int n, bool inv = false)
                                                                 static inline NTT<MAXN, P, RT> ntt;
                                                                 int n() const { return (int)size(); } // n() >= 1
    bitrev(a, n);
                                                                 Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
                                                                      .data(), min(p.n(), m), data()); }
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
                                                                 Poly &irev() { return reverse(data(), data() + n()),
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + dl; ++j, x += dx
                                                                 Poly &isz(int m) { return resize(m), V; }
                                                                 static ll minv(ll x) { return ntt.minv(x); }
           val_t + mp = a[j + dl] * (inv ? conj(w[x]) : w
                                                                 == fb1867 ==
                                                                 Poly &iadd(const Poly &rhs) { // db5668
               [x]);
          a[j + dl] = a[j] - tmp;
                                                                   fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
                                                                   return V; // need n() == rhs.n()
          a[j] += tmp;
                                                                 Poly &imul(11 k) { // a8df26
      }
                                                                   fi(0, n()) V[i] = V[i] * k % P;
    if (inv) {
                                                                   return V:
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
                                                                 Poly Mul(const Poly &rhs) const { // 46caf3
                                                                   int m = 1;
  //multiplying two polynomials A * B:
                                                                   while (m < n() + rhs.n() - 1) m <<= 1;</pre>
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
                                                                   assert(m <= MAXN);</pre>
                                                                   Poly X(V, m), Y(rhs, m);
  //A[i] *= B[i], fft.trans(A, siz, 1);
                                                                   ntt(X, m), ntt(Y, m);
                                                                   fi(0, m) X[i] = X[i] * Y[i] % P;
7.3 NTT [bf683f]
                                                                   ntt(X, m, true);
                                                                   return X.isz(n() + rhs.n() - 1);
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
                                                                 Poly Inv() const { // 796a37
                                                                   if (n() == 1) return {minv(V[0])};
//51*(2^25)+1, 1711276033, 29
                                                                   int m = 1; // need V[0] != 0, 2*sz<=MAXN</pre>
// only works when sz(A) + sz(B) - 1 <= MAXN template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
                                                                   while (m < n() * 2) m <<= 1;</pre>
                                                                   assert(m <= MAXN);</pre>
struct NTT {
                                                                   Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);
  11 w[MAXN];
                                                                   Poly Y(V, m);
  11 mpow(ll a, ll n);
                                                                   ntt(Xi, m), ntt(Y, m);
  11 minv(ll a) { return mpow(a, P - 2); }
                                                                   fi(0, m) {
  NTT() {
                                                                     Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    ll dw = mpow(RT, (P - 1) / MAXN);
                                                                     if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
                                                                   ntt(Xi, m, true);
          % P;
                                                                   return Xi.isz(n());
  void bitrev(vector<11> &a, int n) {
                                                                 Poly &shift_inplace(const 11 &c) { // 0c04f6
    int i = 0;
                                                                   int n = V.n(); // 2 * sz <= MAXN</pre>
    for (int j = 1; j < n - 1; ++j) {</pre>
                                                                   vector<ll> fc(n), ifc(n);
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
                                                                   fc[0] = ifc[0] = 1;
      if (j < i) swap(a[i], a[j]);</pre>
                                                                   for (int i = 1; i < n; i++) {
  fc[i] = fc[i - 1] * i % P;</pre>
    }
                                                                     ifc[i] = minv(fc[i]);
  void operator()(vector<ll> &a, int n, bool inv =
      false) { //0 <= a[i] < P
                                                                   for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
                                                                   Poly g(n);
      int dx = MAXN / L, dl = L >> 1;
                                                                   11 cp = 1;
      for (int i = 0; i < n; i += L) {</pre>
                                                                   for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
                                                                         cp = cp * c % P;
             ) {
                                                                   V = V.irev().Mul(g).isz(n).irev();
           ll tmp = a[j + dl] * w[x] % P;
                                                                   for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
           if ((a[j + d1] = a[j] - tmp) < 0) a[j + d1]
                                                                       Ρ;
                                                                   return V;
          if ((a[j] += tmp) >= P) a[j] -= P;
                                                               // == 7b2835 ==
      }
                                                                 Poly shift(const 11 &c) const { return Poly(V).
                                                                     shift_inplace(c); }
    if (inv) {
                                                                 Poly _Sqrt() const { // Jacobi(V[0], P) = 1
  if (n() == 1) return {QuadraticResidue(V[0], P)};
      reverse(a.begin()+1, a.begin()+n);
      11 invn = minv(n);
                                                                   Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
                                                                   return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
          Ρ;
  }
                                                                 == b46641 ==
};
                                                                 Poly Sqrt() const { // 1aa942
                                                                   Poly a; // 2 * sz <= MAXN
      Polynomial Operation [77a8a8]
                                                                   bool has = 0;
                                                                   for (int i = 0; i < n(); i++) {</pre>
// == b4233a ==
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
                                                                     if (V[i]) has = 1;
```

if (has) a.push_back(V[i]);

```
if (!has) return V;
    if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
      return Poly();
    a = a.isz((n() + a.n()) / 2)._Sqrt();
    int sz = a.n();
    a.isz(n());
    rotate(a.begin(), a.begin() + sz, a.end());
    return a:
  pair<Poly, Poly> DivMod(const Poly &rhs) const { // 5
      bd174
    if (n() < rhs.n()) return {{0}, V};</pre>
    const int m = n() - rhs.n() + 1;
    Poly X(rhs); // (rhs.)back() != 0
    X.irev().isz(m);
    Poly Y(V);
    Y.irev().isz(m);
    Poly Q = Y.Mul(X.Inv()).isz(m).irev();
    X = rhs.Mul(Q), Y = V;
    fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
    return {Q, Y.isz(max(1, rhs.n() - 1))};
  }
// == 76b1af ==
 Poly Dx() const {
    Poly ret(n() - 1);
    fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
    return ret.isz(max(1, ret.n()));
  Poly Sx() const {
    Poly ret(n() + 1);
    fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
    return ret;
  Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
    return Poly(Y.data() + n() - 1, Y.data() + Y.n());
// == 3afa3f ==
  vector<ll> _eval(const vector<ll> &x, const vector<</pre>
      Poly> &up) const { // fb6553
    const int m = (int)x.size();
    if (!m) return {};
    vector<Poly> down(m * 2);
    // down[1] = DivMod(up[1]).second;
    // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
         .second:
    down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
        _tmul(m, V);
    fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
        1, down[i / 2]);
    vector<11> y(m);
    fi(0, m) y[i] = down[m + i][0];
  static vector<Poly> _tree1(const vector<ll> &x) { //
      f5c433
    const int m = (int)x.size();
    vector<Poly> up(m * 2);
    fi(0, m) up[m + i] = {neg(x[i]), 1};
for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
        Mul(up[i * 2 + 1]);
    return up;
  vector<ll> Eval(const vector<ll> &x) const { // 1e5,
    auto up = _tree1(x);
    return _eval(x, up);
  static Poly Interpolate(const vector<11> &x, const
     vector<ll> &y) { // d7bae4
    const int m = (int)x.size(); // 1e5, 1.4s
    vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
    fi(0, m) z[i] = y[i] * minv(z[i]) % P;
    fi(0, m) down[m + i] = {z[i]};
    for (int i = m - 1; i > 0; --i)
      down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
          down[i * 2 + 1].Mul(up[i * 2]));
    return down[1];
```

```
// == c066ab ==
  Poly Ln() const \{ // V[\theta] == 1, 2*sz <= MAXN \}
    return Dx().Mul(Inv()).Sx().isz(n());
  Poly Exp() const \{ // V[0] == 0,2*sz <= MAXN \}
    if (n() == 1) return {1};
    Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
    Poly Y = X.Ln();
    Y[0] = P - 1;
    fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
    return X.Mul(Y).isz(n());
// == 3f1d86 ==
  // M := P(P - 1). If k >= M, k := k % M + M.
  Poly Pow(11 k) const { // 2*sz<=MAXN // d08261
    int nz = 0;
    while (nz < n() && !V[nz]) ++nz;</pre>
    if (nz * min(k, (11)n()) >= n()) return Poly(n());
if (!k) return Poly(Poly{1}, n());
    Poly X(data() + nz, data() + nz + n() - nz * k);
    const 11 c = ntt.mpow(X[0], k % (P - 1));
    return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
         n()).irev();
  // sum_j w_j [x^j] f(x^i) for i \in [0, m]
  Poly power_projection(Poly wt, int m) { // 277119
    assert(n() == wt.n()); // 4*sz <= MAXN!
    if (!n()) {
      return Poly(m + 1, 0);
    if (V[0] != 0) {
      11 c = V[0];
      V[0] = 0;
      Poly A = V.power_projection(wt, m);
      fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
          factorial
      Poly B(m + 1);
      11 pow = 1;
      fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
            * c % P; // inv. of fac
      A = A.Mul(B).isz(m + 1);
      fi(0, m + 1) A[i] = A[i] * fac[i] % P;
      return A;
    int n = 1;
    while (n < V.n()) n *= 2;</pre>
    isz(n), wt.isz(n).irev();
    int k = 1;
    Poly p(wt, 2 * n), q(V, 2 * n);
    q.imul(P - 1);
    while (n > 1) {
      Poly r(2 * n * k);
      fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
          q[i]));
      Poly pq = p.Mul(r).isz(4 * n * k);
      Poly qq = q.Mul(r).isz(4 * n * k);
      fi(0, 2 * n * k) {
        pq[2 * n * k + i] += p[i];
        qq[2 * n * k + i] += q[i] + r[i];
pq[2 * n * k + i] %= P;
         qq[2 * n * k + i] %= P;
      fill(p.begin(), p.end(), 0);
      fill(q.begin(), q.end(), 0);
      for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
        p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];

q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];
      n /= 2, k *= 2;
    Poly ans(k);
    fi(0, k) ans[i] = p[2 * i];
    return ans.irev().isz(m + 1);
  Poly FPSinv() { // 2c54b4

const int n = V.n() - 1;
    if (n == -1) return {};
    assert(V[0] == 0);
    if (n == 0) return V;
    assert(V[1] != 0);
```

```
ll c = V[1], ic = minv(c);
    imul(ic);
    Poly wt(n + 1);
    wt[n] = 1;
    Poly A = V.power_projection(wt, n);
    Poly g(n);
    fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
    g = g.Pow(neg(minv(n)));
    g.insert(g.begin(), 0);
    fi(0, g.n()) g[i] = g[i] * pow % P, pow = pow * ic
         % P;
    return g;
  Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
       1 = k; // 7b552c
    return Poly(*this).irev().Mul(rhs).isz(n()).irev();
  Poly FPScomp(Poly g) { // solves V(g(x)) // 332bb2
    auto rec = [&](auto &rec, int n, int k, Poly Q) ->
         Poly {
       if (n == 1) {
         Poly p(2 * k);
         irev();
         fi(0, k) p[2 * i] = V[i];
         return p;
      Poly R(2 * n * k);
       fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
           Q[i]));
       Poly QQ = Q.Mul(R).isz(4 * n * k);
       fi(0, 2 * n * k) {
         QQ[2 * n * k + i] += Q[i] + R[i];
         QQ[2 * n * k + i] %= P;
      Poly nxt_Q(2 * n * k);
       for(int \bar{j} = 0; j < 2 * k; j++) fi(0, n / 2) {
         nxt_{Q[n * j + i]} = QQ[(2 * n) * j + (2 * i + 0)]
      Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
       Poly pq(4 * n * k);
       for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
   pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +</pre>
         pq[(2*n)*j+(2*i+1)] \%= P;
      Poly p(2 * n * k);
      fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i
           ]) % P;
       pq.pop_back();
      Poly x = pq.TMul(R);
fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
      return p;
    };
    int sz = 1;
    while(sz < n() || sz < g.n()) sz <<= 1;
    return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
          sz)).isz(sz).irev();
  }
};
#undef fi
#undef V
using Poly_t = Poly<1 << 19, 998244353, 3>;
7.5 Generating Function
Ordinary Generating Function
• C(x) = A(rx): c_n = r^n a_n 的一般生成函數。
• C(x) = A(x) + B(x): c_n = a_n + b_n 的一般生成函數。
• C(x) = A(x)B(x): c_n = \sum_{i=0}^n a_i b_{n-i} 的一般生成函數。
• C(x) = A(x)^k: c_n = \sum_{i_1+i_2+\dots+i_k=n} a_{i_1}a_{i_2}\dots a_{i_k} 的一般生成函數。
• C(x) = xA(x)': c_n = na_n 的一般生成函數。
• C(x) = \frac{A(x)}{1-x}: c_n = \sum_{i=0}^n a_i 的一般生成函數。
• C(x) = A(1) + x \frac{A(1) - A(x)}{1 - x}: c_n = \sum_{i=n}^{\infty} a_i 的一般生成函數。
```

常用展開式

•
$$\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$$

• $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$, ${a \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

• 卡特蘭數: $f(x) = \frac{1-\sqrt{1-4x}}{2}$

Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數 $\hat{C}(x) = \hat{A}(x)^k$: $\sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1i_2,\dots,i_k} a_i a_{i_2}\dots a_{i_k}$ 的指數生成函數
- $\hat{C}(x) = \exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 F(0),G(0)=0, $F'(0),G'(0)\neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$

$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

Bostan Mori [41c3bc] 7.6

```
const 11 mod = 998244353;
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = {P, Q};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  vector<11> Qn(sz);
  vector N(2, vector<ll>(sz));
  while(k) {
    fill(iter(Qn), 0);
for(int i = 0; i < d + 1; i++){
      Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    ntt(Qn, sz, false);
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){
      fill(iter(N[i]), 0);
      copy(iter(M[i]), N[i].begin());
      ntt(N[i], sz, false);
for(int j = 0; j < sz; j++)</pre>
        N[i][j] = N[i][j] * Qn[j] % mod;
      ntt(N[i], sz, true);
      for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
        M[i][j >> 1] = N[i][j];
    k \gg 1;
  return M[0][0] * ntt.minv(M[1][0]) % mod;
11 LinearRecursion(vector<ll> a, vector<ll> c, ll k) {
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
  auto A = a; A.resize(sz);
  auto C = c; C.resize(sz);
  ntt(A, sz, false), ntt(C, sz, false);
```

```
for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
                                                                                            // sa[0] = s.size()
  ntt(A, sz, true);
                                                                                            // character should be 1-
  A.resize(d);
                                                                                                 based
                                                               SuffixArray(string& s, int lim=256) { // or
  return BostanMori(A, c, k);
                                                                    basic_string<int>
}
                                                                  int n = s.size() + 1, k = 0, a, b;
                                                                  vector<int> x(n, 0), y(n), ws(max(n, lim));
8
     String
                                                                  rank.assign(n, 0);
                                                                  for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
      KMP Algorithm [c8b75f]
                                                                 sa = lcp = y, iota(sa.begin(), sa.end(), 0);
for (int j = 0, p = 0; p < n; j = max(1, j * 2),</pre>
// 0-based
                                                                      lim = p) {
// fail[i] = max k<i s.t. s[0..k] = s[i-k..i]
                                                                    p = j, iota(y.begin(), y.end(), n - j);
vector<int> kmp_build_fail(const string &s){
                                                                    for (int i = 0; i < n; i++)</pre>
  int n = SZ(s);
                                                                      if (sa[i] >= j) y[p++] = sa[i] - j;
  vector<int> fail(n, -1);
                                                                    for (int &i : ws) i = 0;
  int cur = -1;
                                                                    for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
  for(int i = 1; i < n; i++){</pre>
                                                                    for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];</pre>
    while(cur != -1 && s[cur + 1] != s[i])
                                                                    for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
      cur = fail[cur];
                                                                    swap(x, y), p = 1, x[sa[0]] = 0;
    if(s[cur + 1] == s[i])
                                                                    for(int i = 1; i < n; i++){</pre>
      cur++;
                                                                      a = sa[i - 1], b = sa[i];
    fail[i] = cur;
                                                                      x[b] = (y[a] == y[b] && y[a + j] == y[b + j])?
                                                                           p - 1 : p++:
  return fail;
                                                                   }
                                                                 }
                                                                  for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
void kmp_match(const string &s, const vector<int> &fail
                                                                  for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
    , const string &t){
                                                                    for (k \&\& k--, j = sa[rank[i] - 1];
  int cur = -1;
                                                                        s[i + k] == s[j + k]; k++);
  int n = SZ(s), m = SZ(t);
                                                               }
  for(int i = 0; i < m; i++){</pre>
                                                             };
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t
                                                             8.5 Suffix Automaton [016373]
        [i]))
      cur = fail[cur];
                                                             struct exSAM {
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
                                                               const int CNUM = 26;
      cur++;
                                                                // len: maxlength, link: fail link
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
                                                               // LenSorted: topo order, cnt: occur
  }
                                                               vector<int> len, link, lenSorted, cnt;
                                                               vector<vector<int>> next;
8.2 Manacher Algorithm [caf0f4]
                                                               int total = 0;
                                                               int newnode() {
/* center i: radius z[i * 2 + 1] / 2
                                                                  return total++;
   center i, i + 1: radius z[i * 2 + 2] / 2
   both aba, abba have radius 2 */
                                                               void init(int n) { // total number of characters
len.assign(2 * n, 0); link.assign(2 * n, 0);
vector<int> manacher(const string &tmp){ // 0-based
                                                                  lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
  string s = "%";
  int 1 = 0, r = 0;
                                                                  next.assign(2 * n, vector<int>(CNUM));
  for(char c : tmp) s += c, s += '%';
                                                                  newnode(), link[0] = -1;
  vector<int> z(SZ(s));
  for(int i = 0; i < SZ(s); i++){</pre>
                                                               int insertSAM(int last, int c) {
    z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
                                                                  // not exSAM: cur = newnode(), p = Last
    while(i - z[i] >= 0 \&\& i + z[i] < SZ(s)
                                                                  int cur = next[last][c];
           && s[i + z[i]] == s[i - z[i]])
                                                                  len[cur] = len[last] + 1;
                                                                  int p = link[last];
      ++z[i];
    if(z[i] + i > r) r = z[i] + i, l = i;
                                                                  while (p != -1 && !next[p][c])
  }
                                                                    next[p][c] = cur, p = link[p];
  return z;
                                                                  if (p == -1) return link[cur] = 0, cur;
}
                                                                  int q = next[p][c];
                                                                  if (len[p] + 1 == len[q]) return link[cur] = q, cur
8.3 Lyndon Factorization [7c612b]
                                                                  int clone = newnode();
// partition s = w[0] + w[1] + ... + w[k-1],
                                                                  for (int i = 0; i < CNUM; ++i)
// w[0] >= w[1] >= ... >= w[k-1]
                                                                    next[clone][i] = len[next[q][i]] ? next[q][i] :
// each w[i] strictly smaller than all its suffix
void duval(const string &s, vector<pii> &w) {
                                                                  len[clone] = len[p] + 1;
  for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
                                                                  while (p != -1 && next[p][c] == q)
    for (j = i + 1, k = i; j < n \&\& s[k] <= s[j]; j++)
                                                                    next[p][c] = clone, p = link[p];
      k = (s[k] < s[j] ? i : k + 1);
                                                                  link[link[cur] = clone] = link[q];
    // if (i < n / 2 && j >= n / 2) {
                                                                  link[q] = clone;
    // for min cyclic shift, call duval(s + s)
                                                                 return cur;
    // then here s.substr(i, n / 2) is min cyclic shift
                                                               void insert(const string &s) {
    for (; i <= k; i += j - k)
                                                                  int cur = 0;
      w.pb(pii(i, j - k)); // s.substr(l, len)
                                                                  for (auto ch : s) {
  }
                                                                    int &nxt = next[cur][int(ch - 'a')];
}
                                                                    if (!nxt) nxt = newnode();
8.4 Suffix Array [cd67ea]
                                                                    cnt[cur = nxt] += 1;
                                                                 }
struct SuffixArray {
```

void build() {

queue<int> q;

vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]

and sa[i-1]

```
q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    }
    vector<int> lc(total);
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
        i]]] = i;
  void solve() {
    for (int i = total - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
  }
};
8.6 Z-value Algorithm [488d87]
// z[i] = max \ k \ s.t. \ s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z function(const string &s){
  int n = s.size();
  vector<int> z(n);
  for(int i = 1, l = 0, r = 0; i < n; i++){</pre>
    if(i <= r) z[i] = min(r - i + 1, z[i - 1]);</pre>
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
     z[i]++;
    if(i + z[i] - 1 > r)
      l = i, r = i + z[i] - 1;
  }
  return z;
}
8.7 Main Lorentz [fcfb8f]
struct Rep{ int minl, maxl, len; };
vector<Rep> rep; // 0-base
// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)
void main_lorentz(const string &s, int sft = 0) {
  const int n = s.size();
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
  const string u = s.substr(0, nu), v = s.substr(nu),
        ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
            ());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const auto z1 = z_function(ru), z2 = z_function(v + '
      #' + u),
             z3 = z_{function}(ru + '#' + rv), z4 =
                  z_function(v);
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
  auto add_rep = [&](bool left, int c, int l, int k1,
      int k2) {
    const int L = max(1, 1 - k2), R = min(1 - left, k1)
    if (L > R) return;
    if (left) rep.emplace_back(Rep({sft + c - R, sft +
        c - L, 1}));
    else rep.emplace_back(Rep({sft + c - R - l + 1, sft
         + c - L - 1 + 1, 1));
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int 1, k1, k2;
    if (cntr < nu) {</pre>
      1 = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
}
```

```
8.8 AC Automaton [f529e6]
```

```
const int SIGMA = 26;
struct AC_Automaton {
  // child: trie, next: automaton
  vector<vector<int>> child, next;
  vector<int> fail, cnt, ord;
  int total = 0;
  int newnode() {
    return total++;
  void init(int len) { // len >= 1 + total len
    child.assign(len, vector<int>(26, -1));
    next.assign(len, vector<int>(26, -1));
    fail.assign(len, -1); cnt.assign(len, 0);
    ord.clear();
    newnode();
  int input(string &s) {
    int cur = 0;
    for (char c : s) {
      if (child[cur][c - 'A'] == -1)
    child[cur][c - 'A'] = newnode();
      cur = child[cur][c - 'A'];
    return cur; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(0), fail[0] = -1;
    while(!q.empty()) {
      int R = q.front();
      q.pop(); ord.pb(R);
      for (int i = 0; i < SIGMA; i++)</pre>
        if (child[R][i] != -1) {
          int X = next[R][i] = child[R][i], Z = fail[R
          while (Z != -1 && child[Z][i] == -1)
            Z = fail[Z];
          fail[X] = Z != -1 ? child[Z][i] : 0;
          q.push(X);
        else next[R][i] = R ? next[fail[R]][i] : 0;
    }
  void solve() {
    for (int i : ord | views::reverse)
      if (i) cnt[fail[i]] += cnt[i];
  }
};
      Palindrome Automaton [8a071b]
```

8.9

```
struct PalindromicTree {
 struct node {
    int nxt[26], fail, len; // num = depth of fail link
    int cnt, num; // cnt = occur, num = #pal_suffix of
        this node
    node(int l = 0) : nxt{},fail(0),len(1),cnt(0),num
        (0) {}
 vector<node> st; vector<int> s; int last, n;
 void init() {
    st.clear(); s.clear(); last = 1; n = 0;
    st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
 int getFail(int x) {
    while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x;
 void add(int c) {
   s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
      int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
      st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
```

```
}
void dpcnt() {
  for(int i = SZ(st) - 1; i >= 0; i--){
    auto nd = st[i];
    st[nd.fail].cnt += nd.cnt;
  }
}
int size() { return (int)st.size() - 2; }
;;
```

9 Misc

9.1 Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

9.2 Matroid

 $M=(E,\mathcal{I})$, where $\mathcal{I}\subseteq 2^E$ is nonempty, is a matroid if:

- If $S \in \mathcal{I}$ and $S' \subsetneq S$, then $S' \in \mathcal{I}$.
- For $S_1,S_2\in\mathcal{I}$ s.t. $|S_1|<|S_2|$, there exists $e\in S_2\setminus S_1$ s.t. $S_1\cup\{e\}\in\mathcal{I}$. Matroid intersection:

Start from $S = \emptyset$. In each iteration, let

- $Y_1 = \{x \notin S \mid S \cup \{x\} \in \mathcal{I}_1\}$
- $Y_2 = \{x \not\in S \mid S \cup \{x\} \in \mathcal{I}_2\}$

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \notin S$, create edges

- $x \to y$ if $S \{x\} \cup \{y\} \in \mathcal{I}_1$.
- $y \to x$ if $S \{x\} \cup \{y\} \in \mathcal{I}_2$.

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.3 Simulate Annealing [ff826c]

```
ld anneal() {
 mt19937 rnd_engine(seed);
  uniform_real_distribution<ld> rnd(0, 1);
  const ld dT = 0.001;
  // Argument p
 ld S_cur = calc(p), S_best = S_cur;
  for (ld T = 2000; T > eps; T -= dT) {
   // Modify p to p_prime
    const ld S_prime = calc(p_prime);
    const ld delta_c = S_prime - S_cur;
    ld prob = min((ld)1, exp(-delta_c / T));
    if (rnd(rnd_engine) <= prob)</pre>
      S_cur = S_prime, p = p_prime;
    if (S_prime < S_best) // find min</pre>
      S_best = S_prime, p_best = p_prime;
  return S_best;
```

9.4 Binary Search On Fraction [f6b9ec]

```
struct Q {
    11 p, q;
    Q go(Q b, ll d) { return {p + b.p * d, q + b.q * d};
      }
};

// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N, auto &&pred) {
    Q lo{0, 1}, hi{1, 0};
    if (pred(lo)) return lo;
    assert(pred(hi));
    bool dir = 1, L = 1, H = 1;
    for (; L || H; dir = !dir) {
      ll len = 0, step = 1;
}</pre>
```

```
for (int t = 0; t < 2 && (t ? step /= 2 : step *=</pre>
        2);)
      if (Q mid = hi.go(lo, len + step);
          mid.p > N \mid\mid mid.q > N \mid\mid dir \land pred(mid))
        t++:
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
9.5 Min Plus Convolution [09b5c3]
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector
    <int> &b) +
  int n = SZ(a), m = SZ(b);
  vector<int> c(n + m - 1, INF);
  auto dc = [&](auto Y, int 1, int r, int jl, int jr) {
    if (1 > r) return;
    int mid = (1 + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)
  if (int i = mid - j; i >= 0 && i < n)</pre>
        if (best > a[i] + b[j])
          best = a[i] + b[j], from = j;
    Y(Y, l, mid - 1, jl, from), Y(Y, mid + 1, r, from,
        jr);
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
9.6 SMAWK [a2a4ce]
// For all 2x2 submatrix:
// If M[1][0] < M[1][1], M[0][0] < M[0][1]
// If M[1][0] == M[1][1], M[0][0] <= M[0][1]
// M[i][ans_i] is the best value in the i-th row
// select(int r, int u, int v) return true if f(r, v)
    is better than f(r, u)
vector<int> smawk(int N, int M, auto &&select) {
  auto dc = [&](auto self, const vector<int> &r, const
      vector<int> &c) {
    if (r.empty()) return vector<int>{};
    const int n = SZ(r); vector<int> ans(n), nr, nc;
    for (int i : c) {
      while (!nc.empty() &&
          select(r[nc.size() - 1], nc.back(), i))
        nc.pop_back();
      if (int(nc.size()) < n) nc.push_back(i);</pre>
    for (int i = 1; i < n; i += 2) nr.push_back(r[i]);</pre>
    const auto na = self(self, nr, nc);
    for (int i = 1; i < n; i += 2) ans[i] = na[i >> 1];
    for (int i = 0, j = 0; i < n; i += 2) {</pre>
      ans[i] = nc[j];
      const int end = i + 1 == n ? nc.back() : ans[i +
      while (nc[j] != end)
        if (select(r[i], ans[i], nc[++j])) ans[i] = nc[
             j];
    return ans;
  vector<int> R(N), C(M); iota(iter(R), 0), iota(iter(C)
      ), 0);
  return dc(dc, R, C);
9.7 Golden Ratio Search [ce06a8]
ld goldenRatioSearch(ld a, ld b, auto &&f) {
  1d r = (sqrt(5)-1)/2, eps = 1e-7;
  1d x1 = b - r*(b-a), x2 = a + r*(b-a);
  ld f1 = f(x1), f2 = f(x2);
  while (b-a > eps)
    if (f1 < f2) { //change to > to find maximum
      b = x2; x2 = x1; f2 = f1;
      x1 = b - r*(b-a); f1 = f(x1);
    } else {
      a = x1; x1 = x2; f1 = f2;
      x2 = a + r*(b-a); f2 = f(x2);
  return a;
```

10 Notes

10.1 Geometry

Rotation Matrix

$$\begin{pmatrix}
\cos\theta & -\sin\theta \\
\sin\theta & \cos\theta
\end{pmatrix}$$

- rotate 90° : $(x,y) \rightarrow (-y,x)$
- rotate -90° : $(x,y) \rightarrow (y,-x)$

Triangles

Side lengths: a, b, c

Semiperimeter:
$$p = \frac{a+b+c}{2}$$

Semiperimeter:
$$p=\frac{a+b+c}{2}$$
 Area: $A=\sqrt{p(p-a)(p-b)(p-c)}$ Circumradius: $R=\frac{abc}{4A}$

Circumradius:
$$R = \frac{abc}{}$$

Inradius:
$$r = \frac{A}{n}$$

Length of median (divides triangle into two equal-area triangles): $m_a=rac{1}{2}\sqrt{2b^2+2c^2-a^2}$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}$$

Law of sines: $\frac{\sin\alpha}{a}=\frac{\sin\beta}{b}=\frac{\sin\gamma}{c}=\frac{1}{2R}$ Law of cosines: $a^2=b^2+c^2-2bc\cos\alpha$

Law of cosines:
$$a^2 = b^2 + c^2 - 2bc$$
 Co
$$\frac{a+b}{a-b} = \frac{\tan\frac{\alpha+\beta}{2}}{\tan\frac{\alpha-\beta}{2}}$$
Ouadrilaterals

Ouadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A=2ef\cdot\sin\theta=F\tan\theta=\sqrt{4e^2f^2-F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef=ac+bd , and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \mathrm{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \mathrm{atan2}(y,x) \end{array}$$

Green's Theorem

$$\begin{split} \iint_D \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy &= \oint_{L^+} (P dx + Q dy) \\ \text{Area} &= \frac{1}{2} \oint_{\Gamma} x \, dy - y \, dx \end{split}$$

· Circular sector:

$$\begin{split} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r(r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{split}$$

· Centroid:

$$\bar{x} = \frac{1}{2A} \int_C x^2 dy \quad \bar{y} = -\frac{1}{2A} \int_C y^2 dx$$

Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1, p_2, p_3 are collinear $\iff p_1^*, p_2^*, p_3^*$ intersect at a point
- p lies above $l \iff l^*$ lies above p^*
- lower convex hull \leftrightarrow upper envelope

10.2 Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x}) \qquad \qquad \cosh x = \frac{1}{2}(e^x + e^{-x}) \\ \sin n\pi = 0 \qquad \qquad \cos n\pi = (-1)^n$$

$$\begin{aligned} \sin(\alpha+\beta) &= \sin\alpha\cos\beta + \cos\alpha\sin\beta\\ \cos(\alpha+\beta) &= \cos\alpha\cos\beta - \sin\alpha\sin\beta\\ \sin(2\alpha) &= 2\cos\alpha\sin\alpha\\ \cos(2\alpha) &= \cos^2\alpha - \sin^2\alpha = 2\cos^2\alpha - 1 = 1 - 2\sin^2\alpha \end{aligned}$$

$$\tan(\alpha+\beta) = \frac{\tan\alpha + \tan\beta}{1-\tan\alpha\tan\beta}$$

$$\sin\alpha + \sin\beta = 2\sin\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2}$$

$$\cos\alpha + \cos\beta = 2\cos\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2}$$

$$\sin\alpha\sin\beta = \frac{1}{2}(\cos(\alpha-\beta) - \cos(\alpha+\beta))$$

$$\sin\alpha\cos\beta = \frac{1}{2}(\sin(\alpha+\beta) + \sin(\alpha-\beta))$$

$$\cos\alpha\sin\beta = \frac{1}{2}(\sin(\alpha+\beta) - \sin(\alpha-\beta))$$

$$\cos\alpha\cos\beta = \frac{1}{2}(\cos(\alpha-\beta) + \cos(\alpha+\beta))$$

$$(V+W)\tan(\alpha-\beta)/2 = (V-W)\tan(\alpha+\beta)/2$$
 where V,W are lengths of sides opposite angles α,β .
$$a\cos x + b\sin x = r\cos(x-\phi)$$

$$a\sin x + b\cos x = r\sin(x+\phi)$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

10.3 Calculus

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1-x^{2}}} \qquad \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^{2}}}$$

$$\frac{d}{dx} \tan x = 1 + \tan^{2}x \qquad \frac{d}{dx} \arctan x = \frac{1}{1+x^{2}}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^{2}}$$

$$\int e^{-x^{2}} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int xe^{ax} = \frac{e^{ax}}{a^{2}} (ax - 1)$$

$$\int \sin^{2}(x) = \frac{x}{2} - \frac{1}{4} \sin 2x \qquad \int \sin^{3}x = \frac{1}{12} \cos 3x - \frac{3}{4} \cos x$$

$$\int \cos^{2}(x) = \frac{x}{2} + \frac{1}{4} \sin 2x \qquad \int \cos^{3}x = \frac{1}{12} \sin 3x + \frac{3}{4} \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int xe^{x} = e^{x}(x - 1) \qquad \int x^{2}e^{x} = e^{x}(x^{2} - 2x + 2)$$

$$\int x^{2} \sin x = 2x \sin x - (x^{2} - 2) \cos x$$

$$\int x^{2} \cos x = 2x \cos x + (x^{2} - 2) \sin x$$

$$\int e^{x} \sin x = \frac{1}{2}e^{x}(\sin x - \cos x)$$

$$\int e^{x} \cos x = \frac{1}{2}e^{x}(\sin x - \cos x)$$

$$\int xe^{x} \sin x = \frac{1}{2}e^{x}(\sin x - x \cos x - \sin x)$$

$$\int xe^{x} \cos x = \frac{1}{2}e^{x}(x \sin x - x \cos x - \sin x)$$

10.4 Sum & Series

$$\begin{split} c^a + c^{a+1} + \dots + c^b &= \frac{c^{b+1} - c^a}{c - 1}, c \neq 1 \\ 1 + 2 + 3 + \dots + n &= \frac{n(n + 1)}{2} \\ 1^2 + 2^2 + 3^2 + \dots + n^2 &= \frac{n(2n + 1)(n + 1)}{6} \\ 1^3 + 2^3 + 3^3 + \dots + n^3 &= \frac{n^2(n + 1)^2}{4} \\ 1^4 + 2^4 + 3^4 + \dots + n^4 &= \frac{n(n + 1)(2n + 1)(3n^2 + 3n - 1)}{3} \\ e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \ (-\infty < x < \infty) \\ \ln(1 + x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \ (-1 < x \le 1) \\ \sqrt{1 + x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \ (-1 \le x \le 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \ (-\infty < x < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \ (-\infty < x < \infty) \end{split}$$

10.5 Misc

Cramer's rule

$$ax + by = e cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc} y = \frac{af - ec}{ad - bc}$$

• Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii}=d(i)$, $L_{ij}=-c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- BEST theorem: the number of eulerian circuits in a directed graph is $|\det(\tilde{L}_{ww})| \cdot \prod_{v \in V} (\deg(v) - 1)!$.
- Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}.$ $rac{rank(D)}{2}$ is the maximum matching on G.

- Cayley's Formula
 - Given a degree sequence d_1, d_2, \ldots, d_n for each $\emph{labeled}$ vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.
- Erdős-Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on \boldsymbol{n} vertices if and only

if
$$d_1 + \cdots + d_n$$
 is even and $\sum_{i=1}^k d_i \le k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$ holds

for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n

is bigraphic if and only if
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

holds for every 1 < k < n.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$

$$\cdots \geq a_n \text{ is digraphic if and only if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1)$$
Fibonacci numbers:
$$\frac{n}{F_n} \begin{vmatrix} 1 & 2 & 3 & 4 & 5 & 31 & 45 & 88 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1246269 & 1e9 & 1e18 \\ \hline F_n & 1 & 1 & 3 & 5 & 8 & 1e18 \\ \hline F_n & 1 & 1 & 3$$

$$1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

· Pick's theorem

For simple polygon, when points are all integer, we have

 $A=\#\{\mbox{lattice points in the interior}\}+\frac{\#\{\mbox{lattice points on the boundary}\}}{2}-1.$

- Möbius inversion formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
 - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- · Spherical cap
 - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : arcsin(a/r).
 - Volume = $\pi h^2 (3r h)/3 = \pi h(3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \sin \theta)$
 - Area $=2\pi rh=\pi(a^2+h^2)=2\pi r^2(1-\cos\theta).$
- Lagrange multiplier
 - Optimize $f(x_1,\ldots,x_n)$ when k constraints $g_i(x_1,\ldots,x_n)=0$.
 - Lagrangian function

$$\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k) = f(x_1,\ldots,x_n) - \sum_{i=1}^k \lambda_i g_i(x_1,\ldots,x_n)$$

- $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n).$ The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- · Nearest points of two skew lines
 - Line 1 : ${m v}_1 = {m p}_1 + t_1 {m d}_1$
 - Line 2 : $v_2 = p_2 + t_2 d_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
 - $n_1 = d_1 \times n$
 - $\boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n}$

 - $egin{aligned} &m{-} & m{n}_2 = m{d}_2 imes m{n} \ &m{-} & m{c}_1 = m{p}_1 + rac{(m{p}_2 m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1 \ &m{-} & m{c}_2 = m{p}_2 + rac{(m{p}_1 m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2 \end{aligned}$
- Bernoulli numbers

$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\begin{split} &\sum_{j=0}^{m} {m+1 \choose j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}. \\ &S_m(n) = \sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k^+ n^{m+1-k} \end{split}$$

exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

$$\prod_{n=1}^{n} (1-x^n) = 1 + \sum_{k=1}^{n} (-1)^k$$
• Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t.

$$\pi(j) \geq j, k \text{ j:s s.t. } \pi(j) > j.$$

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

10.6 Number

Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039, 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

• Number of paritions of n:

n | 2 3 4 5 6 7 8 9 20 30 40 50 100 p(n) | 2 3 5 7 11 15 22 30 627 5604 4e4 2e5 2e8

Maximum number of divisors:

n |100 1e3 1e6 1e9 1e12 1e15 1e18 d(i) 12 32 240 1344 6720 26880 103680

n 123456789 10 11 12 13 14 15 $\binom{2n}{n}$ 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

- Number of ways to partition a set of \boldsymbol{n} labeled elements:

n | 2 3 4 5 6 7 8 9 10 11 12 13 B_n 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7

