using pii = pair<int, int>;

void debug(){cerr << "\n";}</pre> template < class T, class ... U>

using pll = pair<ll, ll>;

#ifdef zisk

```
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                                                                              size += ptr:
                                                                        }; // remember to call flush
#include <bits/stdc++.h>
using namespace std;
                                                                         1.4 Random [4cf9ed]
#define iter(v) v.begin(),v.end()
#define SZ(v) int(v.size())
                                                                              time_since_epoch().count());
#define pb emplace_back
#define ff first
                                                                         1.5 PBDS Tree [7c702a]
#define ss second
                                                                         #include <bits/extc++.h>
using 11 = long long;
```

```
void debug(T a, U ... b){cerr << a << " ", debug(b...)</pre>
template < class T> void pary(T 1, T r){
ostream& operator<<(ostream& o, pair<A,B> p)
                                 << p.ss << ')'; }
  ios_base::sync_with_stdio(false); cin.tie(0);
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
map <F9> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -
    Wextra -Wshadow -O2 -Dzisk -g -fsanitize=undefined,
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
" -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG
  static char *p = buf , *end = buf;
    if((end = buf + fread(buf , 1 , N , stdin)) == buf)
  char buf[buf_size]; int size = 0, ret;
  void flush() { ret = write(1, buf, size); size = 0; }
  void _flush(int sz) { if (sz + size > buf_size) flush
  void write_char(char c) { _flush(1); buf[size++] = c;
    if (x < 0) buf[size++] = '-', x = -x;</pre>
    if (x == 0) buf[size + (ptr++)] = '0';
    else for (; x; x /= 10) buf[size + (ptr++)] = '0' +
    reverse(buf + size, buf + size + ptr);
mt19937 rng(chrono::system clock::now().
using namespace __gnu_pbds;
using Tree = tree<int, null_type, less<>, rb_tree_tag,
    tree_order_node_statistics_update>;
// .find_by_order(x)
// .order_of_key(x)
```

# 1.6 Pragma [6006f6]

"<style>\*{stroke-width:0.5%;}</style>\n",

void line(auto x1, auto y1, auto x2, auto y2) {
 p("<line x1='\$' y1='\$' x2='\$' y2='\$' stroke='\$'/>\n

void circle(auto x, auto y, auto r) {
 p("<circle cx='\$' cy='\$' r='\$' stroke='\$' "
 "fill='none'/>\n", x, -y, r, c); }
void text(auto x, auto y, string s, int w = 12) {
 "("<tay\* x='\$' y='\$' fort size='\$ny'\\$' (tox\*t) b'</pre>

p("<text x='\$' y='\$' font-size='\$px'>\$</text>\n",

x1, -y2, x2 - x1, y2 - y1); } ~SVG() { p("</svg>\n"); }

void color(string nc) { c = nc; }

 $x1, -y1, x2, -y2, c); }$ 

# 2 Data Structure

};

x, -y, w, s); }

# 2.1 Heavy-Light Decomposition [f2dbca]

```
struct HLD{ // 1-based
  int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
  explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1),
  ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
   }
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  void build(){ dfs(1, 1), cut(1, 1); }
int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // Lca
};
```

# 2.2 Link Cut Tree [cf4f34]

```
// 1-based
// == PART HASH ==
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node() : pa{0}, ch{0, 0}, rev{false}, size{1}, v{},
      sum\{\}, rsum\{\}, sv\{\}, sub\{\}, vir\{\} \{\}
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is_rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false:
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
    cur.rev ^= 1;
// == PART HASH ==
  void rotate(int u) {
    int f = cur.pa, g = o[f].pa, l = is_rch(u);
    if (cur.ch[1 ^ 1]) o[cur.ch[1 ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  void splay(int u) {
    vector<int> stk = {u};
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is_root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    up(u);
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
      cur.vir = cur.vir + o[rc].sub - o[last].sub;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0:
    for (access(u); u; u = lc) down(la = u);
    return la;
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
// == PART HASH ==
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); }
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
```

```
return find_root(u) == find_root(v); }
  void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
#undef rc
};
2.3 Treap [2ac37e]
mt19937 rng(880301);
// == PART HASH ==
struct node {
  11 data; int sz;
  node *1, *r;
  node(11 k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
node pool[1000010]; int pool_cnt = 0;
node *newnode(11 k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
  return b->down(), b->l = merge(a, b->l), b->up(), b;
}
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)</pre>
    a = o, split(o->r, a->r, b, k), a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
// a: size k, b: size n - k
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();</pre>
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
// == PART HASH ==
node *kth(node *o, ll k) { // 1-based
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o\rightarrow 1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, 11 key) { // num of key < key</pre>
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, ll k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, ll k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
```

```
tuple<node*, node*, node*> interval(node *&o, int 1,
    int r) { // 1-based
  node *a, *b, *c; // b: [l, r]
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
  return make_tuple(a, b, c);
}
2.4 KD Tree [375ca2]
namespace kdt {
  int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
  point p[maxn];
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
    function < bool (const point &, const point &) > f =
      [dep](const point &a, const point &b) {
        if (dep & 1) return a.x < b.x;</pre>
        else return a.y < b.y;</pre>
      };
    int m = (1 + r) >> 1;
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
    yl[m] = yr[m] = p[m].y;
    lc[m] = build(1, m, dep + 1);
    if (~lc[m]) {
      x1[m] = min(x1[m], x1[1c[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
    return m;
  bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
        q.y < yl[o] - ds || q.y > yr[o] + ds)
      return false;
    return true;
  long long dist(const point &a, const point &b) {
    return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
  void dfs(
      const point &q, long long &d, int o, int dep = 0)
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x ||</pre>
        !(dep & 1) && q.y < p[o].y) {
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    } else {
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    }
  void init(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
    root = build(0, v.size());
  long long nearest(const point &q) {
    long long res = 1e18;
    dfs(q, res, root);
    return res:
} // namespace kdt
2.5 Leftist Tree [e91538]
struct node {
  11 v, data, sz, sum;
node *1, *r;
  node(ll k)
```

}:

```
: v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a\rightarrow r) \rightarrow V(a\rightarrow l)) swap(a\rightarrow r, a\rightarrow l);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
}
2.6 Convex 1D/1D [a449dd]
```

```
template < class T>
struct DynamicHull {
  struct seg { int x, l, r; };
  T f; int C; deque<seg> dq; // range: 1~C
  explicit DynamicHull(T _f, int _C): f(_f), C(_C) {} // max t s.t. f(x, t) >= f(y, t), x < y, maintain max
  int intersect(int x, int y) {
    int 1 = 0, r = C + 1;
    while (1 + 1 < r) {
      int mid = (1 + r) / 2;
      if (f(x, mid) >= f(y, mid)) l = mid;
      else r = mid;
    }
    return 1;
  void push_back(int x) {
    for (int i; !dq.empty() &&
         (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
         ();
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back(). r + 1 \le C) dq.pb(seg({x, dq.back().}
        r + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
  }
};
```

# Dynamic Convex Hull [b45ebc]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
      rhs.a; }
  bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  bool isect(iterator x, iterator y) {
    if (y == end()) \{ x \rightarrow p = kInf; return 0; \}
    if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf
    else x -> p = iceil(y -> b - x -> b, x -> a - y -> a);
    return x->p >= y->p;
  }
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y =
        erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  11 query(ll x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
```

```
Flow & Matching
3
```

# 3.1 Dinic [801a71]

```
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
    ll to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
    n = _n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
    g[u].pb(edge{v, cap, 0, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
    dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int cur = q.front(); q.pop();
      for (auto &e : g[cur])
        if (dis[e.to] == -1 && e.flow != e.cap) {
          q.push(e.to);
          dis[e.to] = dis[cur] + 1;
        }
      }
    }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
          g[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0:
  11 maxflow(int _s, int _t) {
         _s; t = _t;
    11 \text{ flow} = 0, df;
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
};
3.2
      Bounded Flow [758826]
```

```
struct BoundedFlow : Dinic {
 vector<ll> tot;
 void init(int _n) {
   Dinic::init(_n + 2);
    tot.assign(n, 0);
 void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
```

```
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
                                                             struct MinCostCirculation { // 0-based, O(VE * ElogC)
                                                              struct edge {
  bool feasible() {
                                                                11 from, to, cap, fcap, flow, cost, rev;
    11 \text{ sum} = 0;
    int vs = n - 2, vt = n - 1;
                                                              int n;
    for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
        add_edge(vs, i, 0, tot[i]), sum += tot[i];
      else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
        g[vs].pop_back(), g[i].pop_back();
      else if(tot[i] < 0)</pre>
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
                                                                }:
    add_edge(_t, _s, 0, INF);
    if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    };
                                                                }
                                                              }
3.3 MCMF [671e14]
struct MCMF { // 0-base
  struct Edge {
   11 from, to, cap, flow, cost, rev;
  int n, s, t;
  vector<vector<Edge>> g;
  vector<Edge*> past;
  vector<ll> dis, up, pot;
                                                                  }
  explicit MCMF(int
                     _n): n(_n), g(n), past(n), dis(n),
      up(n), pot(n) {}
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(Edge{a, b, cap, 0, cost, SZ(g[b])});
    g[b].pb(Edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
  bool BellmanFord() {
    vector<bool> inq(n);
    fill(iter(dis), INF);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
                                                                }
        if (!inq[u]) inq[u] = 1, q.push(u);
                                                              }
      }
    };
                                                                n = n:
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
relax(e.to, d2, min(up[u], e.cap - e.flow), &e)
                                                              }
      }
                                                            };
    }
                                                            3.5
    return dis[t] != INF;
  pair<ll, 1l> solve(int _s, int _t, bool neg = true) {
    s = _s, t = _t; 11 flow = 0, cost = 0;
    if (neg) BellmanFord(), pot = dis;
    for (; BellmanFord(); pot = dis) {
      for (int i = 0; i < n; ++i)</pre>
        if (dis[i] != INF) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], g[e.to][e.rev].flow -= up[t];
                                                            }
      }
    return {flow, cost};
};
                                                              int n;
```

# vector<edge\*> past; vector<vector<edge>> g; vector<ll> dis; void BellmanFord(int s) { vector<int> inq(n); dis.assign(n, INF); queue<int> q; auto relax = [&](int u, ll d, edge \*e) { **if** (dis[u] > d) { dis[u] = d, past[u] = e; **if** (!inq[u]) inq[u] = 1, q.push(u); relax(s, 0, 0); while (!q.empty()) { int u = q.front(); q.pop(), inq[u] = 0;for (auto &e : g[u]) if (e.cap > e.flow) relax(e.to, dis[u] + e.cost, &e); void try\_edge(edge &cur) { if (cur.cap > cur.flow) return ++cur.cap, void(); BellmanFord(cur.to); if (dis[cur.from] + cur.cost < 0) {</pre> ++cur.flow, --g[cur.to][cur.rev].flow; for (int i = cur.from; past[i]; i = past[i]->from ) { auto &e = \*past[i]; ++e.flow, --g[e.to][e.rev].flow; ++cur.cap; void solve(int mxlg) { // mxlg >= log(max cap) for (int b = mxlg; b >= 0; --b) { for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) e.cap \*= 2, e.flow \*= 2; for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) if (e.fcap >> b & 1) try\_edge(e); void init(int \_n) { past.assign(n, nullptr); g.assign(n, vector<edge>()); void add\_edge(ll a, ll b, ll cap, ll cost) { $g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a$ == b)}): g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1}); **Gomory Hu** [82d968] void GomoryHu(Dinic &flow) { // 0-based int n = flow.n; vector<int> par(n); for (int i = 1; i < n; ++i) {</pre> flow.reset(); add\_edge(i, par[i], flow.maxflow(i, par[i])); for (int j = i + 1; j < n; ++j)</pre> if (par[j] == par[i] && ~flow.dis[j]) par[j] = i;Stoer Wagner Algorithm [a9917b] struct StoerWagner { // 0-based, 0(V^3) vector<int> vis, del; vector<ll> wei;

#### Min Cost Circulation [47cf18] 3.4

}:

```
vector<vector<ll>> edge;
  void init(int _n) {
   n = _n;
    del.assign(n, 0);
    edge.assign(n, vector<ll>(n));
  void add_edge(int u, int v, ll w) {
   edge[u][v] += w, edge[v][u] += w;
  void search(int &s, int &t) {
   vis.assign(n, 0); wei.assign(n, 0);
    s = t = -1;
    while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
        if (!del[i] && !vis[i] && mx < wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)</pre>
        if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
   }
  ll solve() {
   11 ret = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
      for (int j = 0; j < n; ++j)</pre>
        edge[x][j] = (edge[j][x] += edge[y][j]);
    return ret:
 }
};
3.7
     Bipartite Matching [5bb9be]
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!~my[e] || (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  }
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < n1; ++i)</pre>
     if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
   }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
   n1 = _n1, nr = _nr;
    g.assign(nl, vector<int>());
```

```
Kuhn Munkres Algorithm [683e0a]
```

```
struct KM \{ // O-based, maximum matching, O(V^3) \}
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll> hl, hr, slk;
  vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    // -INF for perfect matching
    w.assign(n, vector<ll>(n, 0));
    pre.assign(n, 0);
    qu.assign(n, 0);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
        0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
               w[x][y])) {
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !s1k[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
         ; hr.assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  3
};
      Max Simple Graph Matching [907d7c]
```

# 3.9

```
struct Matching { // 0-based, O(V^3)
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> g;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
   }
  void Blossom(int x, int y, int 1) {
    for (; Find(x) != 1; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
```

```
}
  bool Bfs(int r) {
    iota(iter(fa), 0); fill(iter(s), -1);
q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
       for (int x = q.front(); int u : g[x])
         if (s[u] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
              for (int a = u, b = x, last;
    b != n; a = last, b = pre[a])
                last = match[b], match[b] = a, match[a] =
                      b;
              return true;
           }
           q.push(match[u]); s[match[u]] = 0;
         } else if (!s[u] && Find(u) != Find(x)) {
           int 1 = LCA(u, x);
           Blossom(x, u, 1); Blossom(u, x, 1);
    return false;
  Matching(\textbf{int} \_n) : n(\_n), fa(n + 1), s(n + 1), vis(n
  + 1), pre(n + 1, n), match(n + 1, n), g(n) {} void add_edge(int u, int v)
  { g[u].pb(v), g[v].pb(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
       if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

# 3.10 Stable Marriage

```
1: Initialize m \in M and w \in \overline{W} to free
 2: while \exists free man m who has a woman w to propose to do
3:
         w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
 4:
         if \exists some pair (m', w) then
              if w prefers m to m' then
 5:
 6:
                  m' \leftarrow \textit{free}
7:
                  (m,w) \leftarrow \mathsf{engaged}
8:
             end if
9:
         else
              (m,w) \leftarrow \textit{engaged}
10:
         end if
11:
```

# 12: end while 3.11 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
- 1. Construct super source S and sink T.
- 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
  - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
  - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
- 2. DFS from unmatched vertices in X.
- 3.  $x \in X$  is chosen iff x is unvisited.
- 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source  ${\cal S}$  and sink  ${\cal T}$
- 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y \to x$  with (cost,cap)=(-c,1)
- c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=(0,d(v))
- 5. For each vertex v with d(v)<0 , connect  $v\to T$  with (cost,cap)=(0,-d(v))

- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source  $s \to v$ ,  $v \in G$  with capacity K
- 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity ...
- 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e)) \cdot 2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. Let  $w'(u,v)=w(u,v)-\mu(u)-\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 2. Find the minimum weight matching M with w' . The answer is  $\sum \mu(v) + w'(M)$ .
- · Project selection problem
  - 1. If  $p_v>0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
  - The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$ , Flow  $f_{uv}$ , Cost  $w_{uv}$ , Required Flow difference for vertex  $h_{vv}$
- 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{aligned} \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} &\geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} &= -b_{u} \end{aligned}$$

# 4 Geometry

# 4.1 Geometry Template [86f0f1]

```
using ld = ll;
using pdd = pair<ld, ld>;
#define X first
#define Y second
// Ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){    return v > eps ? 1 : ( v < -eps ? -1
     : 0);
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)
    ) <= 0; }
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&
  ori(p3, p4, p1) * ori(p3, p4, p2) < 0;</pre>
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
```

```
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(}
    p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) /
     abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
     pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
      dot(dp, dq));
  return q0 + pdd(cross(r - p0, num), dot(r - p0, num))
       / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
```

# 4.2 Polar Angle Comparator [808e89]

# 4.3 Minkowski Sum [98abff]

```
void reorder_poly(vector<pdd>& pnts){
 int mn = 0;
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn</pre>
        ].Y && pnts[i].X < pnts[mn].X))
 rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
  reorder_poly(P);
  reorder_poly(Q);
 int psz = P.size();
  int qsz = Q.size()
 P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
 vector<pdd> ans;
  int i = 0, j = 0;
 while(i < psz || j < qsz){
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1] - P[i], Q[j + 1] - Q[j])
    if(t >= 0) i++;
   if(t <= 0) j++;
  return ans;
}
```

# 4.4 Intersection of Circle and Convex Polygon [63653d]

```
double _area(pdd pa, pdd pb, double r){
 if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
   S = (C/2)*r*r;
   h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
   theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
```

# 4.5 Intersection of Circles [f7a2fe]

# 4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
    int sign1 ){
  vector<Line> ret:
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;

double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
         v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sgn(p1.X - p2.X) == 0 and
         sgn(p1.Y - p2.Y) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret;
```

# **4.7 Intersection of Line and Convex Polygon**[157258]

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 || cmpL(B) > 0)
  return pii(-1, -1); // no collision
auto gao = [&](int 1, int r) {
    for (int t = 1; (1 + 1) % n != r; ) {
      int m = ((1 + r + (1 < r? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(t) ? 1 : r) = m;
    return (1 + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
      +1
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
```

# 4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
    a, pdd b) {
    pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
    ;
    double s = cross(b - a, c - a), h2 = r * r - s * s /
        abs2(b - a);
    if (sgn(h2) < 0) return {};
    if (sgn(h2) == 0) return {p};
    pdd h = (b - a) / abs(b - a) * sqrt(h2);
    return {p - h, p + h};
}</pre>
```

# 4.9 Point in Circle [ecf954]

```
// return q's relation with circumcircle of tri(p[0],p
       [1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
       det += __int128(abs2(p[i]) - abs2(q)) * cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

# 4.10 Point in Convex [f86640]

# 4.11 Half Plane Intersection [dfb833]

```
// from 8BQube
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
     .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(iter(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
      return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
          -1)
      continue:
    while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.
        back()))
      dq.pop_back();
    while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
```

# 4.12 Minimum Enclosing Circle [5af6d5]

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
  random_shuffle(iter(pts));
  pdd c = pts[0];
  1d r = 0;
  for(int i = 1; i < SZ(pts); i++){</pre>
     if(abs(pts[i] - c) <= r) continue;
c = pts[i]; r = 0;</pre>
     for(int j = 0; j < i; j++){
  if(abs(pts[j] - c) <= r) continue;</pre>
       c = (pts[i] + pts[j]) / 2;
       r = abs(pts[i] - c);
        for(int k = 0; k < j; k++){
          if(abs(pts[k] - c) > r)
             tie(c, r) = circumcenter(pts[i], pts[j], pts[
                  k]);
       }
     }
  }
  return {c, r};
```

# 4.13 3D Point [badbbd]

```
// Copy from 8BQube
struct Point {
  double x, y, z;
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
    pi]
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
    .y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
\{ return (a + b + c + d) / 4; \}
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
 Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
```

.x) \* (p3.z - p1.z);

```
Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point rotate_around(Point p, double angle, Point axis)
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
                                                            };
  return u * dot(u, p) * (1 - c) + p * c + cross(u, p)
}
4.14 ConvexHull3D [156311]
struct convex_hull_3D {
struct Face {
  int a. b. c:
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
                                                            struct Edge {
vector<Face> res;
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
                                                              int n;
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn
  (abs2(P[0] - p)) != 0; }));
swap(P[2], *find_if(iter(P), [&](auto p) { return sgn
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(iter(P), [&](auto p) { return sgn
      (volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
    vector<Face> next:
    for (auto f : res) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
    for (auto f : res) {
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
          next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    res = next:
 }
bool same(Face s, Face t) {
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
       return 0;
  return 1;
int polygon_face_num() {
  int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin() + i, [&](
        Face g) { return same(res[i], g); });
  return ans;
double get_volume() {
  double ans = 0;
  for (auto f : res)
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
        ]);
 return fabs(ans / 6);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
                                                              }
      .z) * (p3.y - p1.y);
                                                            };
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
```

```
double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
      .y) * (p3.x - p1.x);
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
       * a + b * b + c * c);
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point
// (0, 0, inf) to avoid degenerate case
4.15 Delaunay Triangulation [6a9916]
/* Delaunay Triangulation:
   Given a sets of points on 2D plane, find a
   triangulation such that no points will strictly
   inside circumcircle of any triangle. */
  int id; // oidx[id]
  list<Edge>::iterator twin;
  Edge(int _id = 0):id(_id) {}
struct Delaunay { // 0-base
  vector<int> oidx;
  vector<list<Edge>> head; // result udir. graph
  vector<pll> p;
  Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
      head(n), p(n) {
    iota(iter(oidx), 0);
    for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(iter(oidx), [&](int a, int b)
        for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
    divide(0, n - 1);
  void addEdge(int u, int v) {
    head[u].push_front(Edge(v));
    head[v].push_front(Edge(u));
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
  void divide(int 1, int r) {
    if (1 == r) return;
    if (1 + 1 == r) return addEdge(1, 1 + 1);
    int mid = (1 + r) >> 1, nw[2] = \{1, r\};
    divide(l, mid), divide(mid + 1, r);
    auto gao = [&](int t) {
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      for (auto it : head[nw[t]]) {
        int v = ori(pt[1], pt[0], p[it.id]);
        if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
            id]) < abs2(pt[1] - pt[0])))
          return nw[t] = it.id, true;
      return false;
    while (gao(0) || gao(1));
    addEdge(nw[0], nw[1]); // add tangent
    while (true) {
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      int ch = -1, sd = 0;
      for (int t = 0; t < 2; ++t)
        for (auto it : head[nw[t]])
          if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
               -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
              id])))
            ch = it.id, sd = t;
      if (ch == -1) break; // upper common tangent
      for (auto it = head[nw[sd]].begin(); it != head[
          nw[sd]].end(); )
        if (seg_strict_intersect(pt[sd], p[it->id], pt[
    sd ^ 1], p[ch]))
          head[it->id].erase(it->twin), head[nw[sd]].
              erase(it++);
        else ++it;
      nw[sd] = ch, addEdge(nw[0], nw[1]);
```

## 4.16 Voronoi Diagram [e4f408]

```
// all coord. is even, you may want to call
    halfPlaneInter after then
                                                             pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
vector<vector<Line>> vec;
void build_voronoi_line(int n, vector<pll> &pts) {
                                                               double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
  Delaunay tool(n, pts); // Delaunay
                                                                     - p2);
  vec.clear(), vec.resize(n);
                                                               double s = a + b + c;
  for (int i = 0; i < n; ++i)
                                                               return (a * p1 + b * p2 + c * p3) / s;
    for (auto e : tool.head[i]) {
      int u = tool.oidx[i], v = tool.oidx[e.id];
                                                             pdd masscenter(pdd p1, pdd p2, pdd p3)
      pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v])
                                                             { return (p1 + p2 + p3) / 3; }
                                                             pdd orthcenter(pdd p1, pdd p2, pdd p3)
          - pts[u]);
      vec[u].pb(Line(m, m + d));
                                                             { return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
    }
                                                                   p3) * 2; }
}
                                                             4.20 Rotating Sweep Line [f5f689]
4.17 Polygon Union [9fbf66]
                                                             struct Event {
 / from 8BQube
                                                               pll d; int u, v;
ld rat(pll a, pll b) {
                                                               bool operator<(const Event &b) const {</pre>
  return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
                                                                 int ret = cmp(d, b.d, false);
  // all poly. should be ccw
                                                                 return ret == -1 ? false : ret; } // no tie-break
ld polyUnion(vector<vector<pll>>> &poly) {
  1d res = 0;
                                                             void rotatingSweepLine(const vector<pll> &p) {
  for (auto &p : poly)
                                                               const int n = SZ(p);
    for (int a = 0; a < SZ(p); ++a) {
                                                               vector<Event> e; e.reserve(n * (n - 1));
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
                                                               for (int i = 0; i < n; i++)</pre>
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
                                                                 for (int j = 0; j < n; j++) // pos[i] < pos[j] when
      for (auto &q : poly) {
                                                                       the event occurs
        if (&p == &q) continue;
                                                                    if (i != j) e.pb(p[j] - p[i], i, j);
        for (int b = 0; b < SZ(q); ++b) {
  pll C = q[b], D = q[(b + 1) % SZ(q)];</pre>
                                                               sort(iter(e));
                                                               vector<int> ord(n), pos(n);
                                                               iota(iter(ord), 0);
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
                                                               sort(iter(ord), [&](int i, int j) { // initial order
            1d sa = cross(D - C, A - C), sb = cross(D - C)
                                                                    return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
                                                                        X < p[j].X; });
                 C, B - C);
            segs.pb(sa / (sa - sb), sgn(sc - sd));
                                                               for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
                                                               // initialize
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D
                                                               for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
                - C)) > 0) {
                                                                 // do something
            segs.pb(rat(C - A, B - A), 1);
                                                                 vector<pii> tmp;
            segs.pb(rat(D - A, B - A), -1);
                                                                 for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
                                                                   tmp.pb(pii(e[j].u, e[j].v));
        }
                                                                  sort(iter(tmp), [&](pii x, pii y){
      }
                                                                      return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
      sort(iter(segs));
                                                                          ], pos[y.ss]); });
                                                                 for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
  tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
                                                                      make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
      for (int j = 1; j < SZ(segs); ++j) {</pre>
                                                                          pos[x]);
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
                                                             }
        cnt += segs[j].Y;
                                                             4.21 Vector In Poly [c6d0fa]
      res += cross(A, B) * sum;
                                                             // ori(a, b, c) >= 0, valid: "strict" angle from a-b to
  return res / 2;
                                                             bool btwangle(pll a, pll b, pll c, pll p, int strict) {
4.18 Tangent Point to Convex Hull [523bc1]
                                                               return ori(a, b, p) >= strict && ori(a, p, c) >=
                                                                    strict;
// from 8BOube
/* The point should be strictly out of hull
                                                             // whether vector{cur, p} in counter-clockwise order
  return arbitrary point on the tangent line */
                                                                 prv, cur, nxt
pii get_tangent(vector<pll> &C, pll p) {
                                                             bool inside(pll prv, pll cur, pll nxt, pll p, int
  auto gao = [&](int s) {
                                                                  strict) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
                                                               if (ori(cur, nxt, prv) >= 0)
    { return ori(p, C[x], C[y]) == s; });
                                                                  return btwangle(cur, nxt, prv, p, strict);
  };
                                                               return !btwangle(cur, prv, nxt, p, !strict);
  return pii(gao(1), gao(-1));
                                                             }
} // return (a, b), ori(p, C[a], C[b]) >= 0
                                                             4.22 Convex Hull DP [92fd4b]
4.19 Heart [082d19]
                                                             sort(iter(pts), [&](pll x, pll y) {
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
                                                                 return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;</pre>
    (center)
                                                                 });
                                                             auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
double m = 2. * (x1 * y2 - y1 * x2);
                                                                  ]; };
                                                             // DP for convex hull vertices (no points on edges)
                                                             auto solve = [\&](int bottom) { // <math>O(n^3)
  pdd center;
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
                                                               pll 0 = pts[bottom];
      y1 - y2)) / m;
                                                               vector<pii> trans;
```

for (int j = bottom + 1; j < n; j++)</pre>

for (int k = bottom + 1; k < n; k++) {</pre>

if (ori(0, pts[j], pts[k]) <= 0) continue;</pre>

center.Y = (x1 \* x2 \* (x2 - x1) - y1 \* y1 \* x2 + x1 \*

y2 \* y2) / m;

return center + p0;

```
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      // check whether j->k is legal
      trans.pb(pii(j, k));
  sort(iter(trans), [&](pii x, pii y) -> bool{
      int tmp = cmp(getvec(x), getvec(y), false);
      if (tmp != -1) return tmp;
      pll v = getvec(x);
      return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
      });
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {</pre>
   // check whether bottom -> j is legal
   // init trans -> j
  for (auto [i, j] : trans) {
   // normal trans i -> j
  for (int j = bottom + 1; j < n; j++) {</pre>
   // check whether j -> bottom is legal
    // end trans j ->
for(int i = 0; i < n; i++) solve(i);</pre>
4.23 Calculate Points in Triangle [bf746f]
// all points are distinct
     i < k < j
  cnt2[i][j] = # of points k s.t. strictly in ij
```

```
// cnt[i][j] = # of point k s.t. strictly above ij, and
// preprocess space: O(n^2), time: O(n^3), query time:
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
  for (int j = 0; j < n; j++){
    if (pts[i] >= pts[j]) continue;
    for (int k = 0; k < n; k++) {
      if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
        int tmp = ori(pts[i], pts[j], pts[k]);
        if (tmp > 0) cnt[i][j]++; // only for i < j</pre>
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
      }
   }
 }
auto calc_tri = [&](array<int, 3> arr) { // strictly
  sort(iter(arr), [\&](int x, int y){return pts[x] <}
      pts[y]; });
  auto [x, y, z] = arr;
  int tmp = ori(pts[x], pts[y], pts[z]);
  if (tmp == 0) return 0;
  else if (tmp < 0)</pre>
    return cnt[x][z] - cnt[x][y] - cnt[y][z] - cnt2[x][
        y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
      x][z];
};
```

### 5 Graph

# 5.1 BCC [d04ebe]

```
struct BCC{ // 0-based, allow multi edges but not allow
     Loops
  int n, m, cnt = 0;
  // n:|V|, m:|E|, cnt:#bcc
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
  vector<vector<int>> bcc_v, bcc_e;
  vector<vector<pii>>> g; // original graph
vector<pii>> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
    n(_n), m(SZ(_edges)), g(_n), edges(_edges){
      for(int i = 0; i < m; i++){</pre>
        auto [u, v] = edges[i];
        g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
  void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
  // modify these if you need more information
  void add v(int v){ bcc v.back().pb(v); }
  void add_e(int e){ bcc_e.back().pb(e); }
  void build(){
    vector < int > in(n, -1), low(n, -1), stk;
```

```
12
    vector<vector<int>> up(n);
    int ts = 0;
    auto _dfs = [&](auto dfs, int now, int par, int pe)
          -> void{
      if(pe != -1) up[now].pb(pe);
      in[now] = low[now] = ts++;
      stk.pb(now);
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
        if(in[v] != -1){
          if(in[v] < in[now]) up[now].pb(e);</pre>
          low[now] = min(low[now], in[v]);
          continue;
        dfs(dfs, v, now, e);
        low[now] = min(low[now], low[v]);
      if((now != par && low[now] >= in[par]) || (now ==
            par && SZ(g[now]) == 0)){
        make_bcc();
        for(int v = stk.back();; v = stk.back()){
           stk.pop_back(), add_v(v);
           for(int e : up[v]) add_e(e);
          if(v == now) break;
        if(now != par) add_v(par);
      }
    };
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) _dfs(_dfs, i, i, -1);
  }
};
5.2 SCC [2c9a01]
struct SCC{ // 0-based, output reversed topo order
  int n, cnt = 0;
  vector<vector<int>> g;
  vector<int> sccid;
  explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
void add_edge(int u, int v){
    g[u].pb(v);
  void build(){
    vector<int> in(n, -1), low(n), stk;
    vector<bool> instk(n);
    int ts = 0;
    auto dfs1 = [&](auto dfs, int now) -> void{
      stk.pb(now); instk[now] = true;
      in[now] = low[now] = ts++;
      for(int i : g[now]){
        if(in[i] == -1)
          dfs(dfs, i), low[now] = min(low[now], low[i])
        else if(instk[i] && in[i] < in[now])</pre>
          low[now] = min(low[now], in[i]);
```

# } };

**5.3 2-SAT** [0686a5]

}

if(low[now] == in[now]){

false;

pop\_back();

for(int i = 0; i < n; i++)</pre>

if(in[i] == -1) dfs1(dfs1, i);

```
struct SAT { // 0-based
  vector<bool> istrue;
  SCC scc;
  SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
  int neg(int a) {
    return a >= n ? a - n : a + n;
  void add_clause(int a, int b) {
    scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
```

for(; stk.back() != now; stk.pop\_back())

sccid[stk.back()] = cnt, instk[stk.back()] =

sccid[now] = cnt++, instk[now] = false, stk.

st[++top] = u;

```
bool solve() {
                                                                                               void reset(int u) {
      scc.build();
                                                                                                  for (int i : vG[u]) reset(i);
      for (int i = 0; i < n; ++i) {</pre>
                                                                                                  vG[u].clear();
         if (scc.sccid[i] == scc.sccid[i + n]) return
                false;
          istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>
                                                                                               void solve(vector<int> &v) {
         istrue[i + n] = !istrue[i];
                                                                                                  top = -1;
                                                                                                  sort(iter(v),
                                                                                                        [&](int a, int b) { return dfn[a] < dfn[b]; });
      return true;
                                                                                                  for (int i : v) insert(i);
   }
                                                                                                  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
};
                                                                                                  // do something
5.4 Dominator Tree [2da9bb]
                                                                                                  reset(vrt);
                                                                                              }
struct Dominator {
   int n;
                                                                                              5.6 Fast DMST [7b274d]
   vector<vector<int>> g, r, rdom; int tk;
vector<int> dfn, rev, fa, sdom, dom, val, rp;
                                                                                               struct E { int s, t; ll w; }; // O-base
   Dominator(int_n): n(n), g(n), r(n), rdom(n), tk(0)
                                                                                               struct PQ {
                                                                                                  struct P {
      dfn = rev = fa = sdom = dom =
                                                                                                     11 v; int i;
         val = rp = vector<int>(n, -1); }
                                                                                                     bool operator>(const P &b) const { return v > b.v;
   void add_edge(int x, int y) { g[x].push_back(y); }
   void dfs(int x) {
                                                                                                  };
      rev[dfn[x] = tk] = x;
                                                                                                  priority_queue<P, vector<P>, greater<>> pq; 11 tag;
      fa[tk] = sdom[tk] = val[tk] = tk; tk++;
                                                                                                         // min heap
      for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
                                                                                                  void push(P p) { p.v -= tag; pq.emplace(p); }
                                                                                                  P top() { P p = pq.top(); p.v += tag; return p; }
                                                                                                  void join(PQ &b) {
         r[dfn[u]].push_back(dfn[x]);
      }
                                                                                                     if (pq.size() < b.pq.size())</pre>
   }
                                                                                                        swap(pq, b.pq), swap(tag, b.tag);
   void merge(int x, int y) { fa[x] = y; }
                                                                                                     while (!b.pq.empty()) push(b.top()), b.pq.pop();
   int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
                                                                                              }; // O(E log^2 V), use leftist tree for O(E log V)
      if (int p = find(fa[x], 1); p != -1) {
                                                                                               vector<int> dmst(const vector<E> &e, int n, int root) {
                                                                                                  vector<PQ> h(n * 2);
for (int i = 0; i < int(e.size()); ++i)</pre>
         if (sdom[val[x]] > sdom[val[fa[x]]])
             val[x] = val[fa[x]];
          fa[x] = p;
                                                                                                     h[e[i].t].push({e[i].w, i});
                                                                                                  vector<int> a(n * 2); iota(iter(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
         return c ? p : val[x];
      } else return c ? fa[x] : val[x];
                                                                                                  auto o = [\&](auto Y, int x) \rightarrow int {
   }
   vector<int> build(int s) {
   // return the father of each node in dominator tree
   dfs(s); // p[i] = -2 if i is unreachable, par[s] =
                                                                                                     return x == a[x] ? x : a[x] = Y(Y, a[x]); };
                                                                                                  auto S = [&](int i) { return o(o, e[i].s); };
                                                                                                  int pc = v[root] = n;
                                                                                                  for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
      for (int i = tk - 1; i >= 0; --i) {
                                                                                                     for (int p = i; v[p]<0 \mid \mid v[p]==i; p = S(r[p])) {
         for (int u : r[i])
                                                                                                        if (v[p] == i)
             sdom[i] = min(sdom[i], sdom[find(u)]);
                                                                                                            for (int q = pc++; p != q; p = S(r[p])) {
          if (i) rdom[sdom[i]].push_back(i);
                                                                                                               h[p].tag -= h[p].top().v; h[q].join(h[p]);
                                                                                                               pa[p] = a[p] = q;
         for (int u : rdom[i]) {
             int p = find(u);
                                                                                                        while (S(h[p].top().i) == p) h[p].pq.pop();
             dom[u] = (sdom[p] == i ? i : p);
                                                                                                        v[p] = i; r[p] = h[p].top().i;
         if (i) merge(i, rp[i]);
      }
                                                                                                  vector<int> ans;
      vector < int > p(n, -2); p[s] = -1;
                                                                                                  for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
      for (int i = 1; i < tk; ++i)</pre>
                                                                                                     for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
         if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
      for (int i = 1; i < tk; ++i)</pre>
                                                                                                        v[f] = n;
         p[rev[i]] = rev[dom[i]];
                                                                                                     ans.push_back(r[i]);
      return p;
   }
                                                                                                  return ans; // default minimize, returns edgeid array
                                                                                              }
};
5.5 Virtual Tree [6abeb5]
                                                                                               5.7 Vizing [58a6ca]
// copy from 8BQube
                                                                                               // find D+1 edge coloring of a graph with max deg D, O(
vector<int> vG[N];
int top, st[N];
                                                                                               struct Vizing { // returns maxdeg+1 edge coloring in
int vrt = -1;
                                                                                                     adjacent matrix G
                                                                                                  int n; // 1-based for vertices and colors, simple
void insert(int u) {
                                                                                                         graph
                                                                                                  vector<vector<int>> C, G;
   if (top == -1) return st[++top] = vrt = u, void();
   int p = LCA(st[top], u);
                                                                                                  vector<int> X, vst;
      if(dep[vrt] > dep[p]) vrt = p;
                                                                                                  Vizing(int _n): n(_n),
   if (p == st[top]) return st[++top] = u, void();
                                                                                                  C(n + 1, vector < int > (n + 2)), G(n + 1, vector < int > (n + 1, 
   while (top >= 1 && dep[st[top - 1]] >= dep[p])
                                                                                                        + 1)),
      vG[st[top - 1]].pb(st[top]), --top;\\
                                                                                                  X(n + 1, 1), vst(n + 1) {}
   if (st[top] != p)
                                                                                                  void solve(vector<pii> &E) {
      vG[p].pb(st[top]), --top, st[++top] = p;
                                                                                                     auto update = [&](int u)
```

{ for (X[u] = 1; C[u][X[u]]; ++X[u]); };

auto color = [&](int u, int v, int c) {

int p = G[u][v];

void dfs(vector<int> &r, vector<int> &c, int 1,

bitset<N> mask) {
while (!r.empty()) {

int p = r.back();

 $r.pop_back()$ , mask[p] = 0;

if (q + c.back() <= ans) return;</pre>

```
G[u][v] = G[v][u] = c;
                                                                     cur[q++] = p;
      C[u][c] = v, C[v][c] = u;
                                                                     vector<int> nr;
      C[u][p] = C[v][p] = 0;
                                                                     for (int i : r) if (G[p][i]) nr.pb(i);
      if (p) X[u] = X[v] = p;
                                                                     if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else update(u), update(v);
                                                                     else if (q > ans) ans = q, copy_n(cur, q, sol);
      return p;
                                                                     c.pop_back(), --q;
    };
    auto flip = [&](int u, int c1, int c2) {
                                                                int solve() {
      int p = C[u][c1];
      swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
                                                                  vector<int> r(n);
                                                                   ans = q = 0, iota(iter(r), 0);
      if (!C[u][c1]) X[u] = c1;
                                                                   pre_dfs(r, 0, bitset<N>(string(n, '1')));
      if (!C[u][c2]) X[u] = c2;
                                                                   return ans:
      return p;
    };
                                                              };
    for (int t = 0; t < SZ(E); ++t) {</pre>
                                                                    Number of Maximal Clique [11fa26]
                                                              5.9
      int u = E[t].ff, v0 = E[t].ss, v = v0, c0 = X[u],
           c = c0, d;
                                                              struct BronKerbosch { // 1-base
      vector<pii> L;
                                                                int n, a[N], g[N][N];
      fill(iter(vst), 0);
                                                                int S, all[N][N], some[N][N], none[N][N];
      while (!G[u][v0]) {
                                                                 void init(int _n) {
        L.emplace_back(v, d = X[v]);
                                                                  n = _n;
for (int i = 1; i <= n; ++i)</pre>
        if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
             --a) c = color(u, L[a].ff, c);
                                                                     for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
        else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
             0; --a) color(u, L[a].ff, L[a].ss);
                                                                 void add_edge(int u, int v) {
        else if (vst[d]) break;
                                                                  g[u][v] = g[v][u] = 1;
        else vst[d] = 1, v = C[u][d];
                                                                void dfs(int d, int an, int sn, int nn) {
      if (!G[u][v0]) {
                                                                  if (S > 1000) return; // pruning
        for (; v; v = flip(v, c, d), swap(c, d));
                                                                   if (sn == 0 && nn == 0) ++S;
        if (int a; C[u][c0]) {
                                                                   int u = some[d][0];
          for (a = SZ(L) - 2; a >= 0 && L[a].ss != c;
                                                                   for (int i = 0; i < sn; ++i) {</pre>
               --a);
                                                                     int v = some[d][i];
          for (; a >= 0; --a) color(u, L[a].ff, L[a].ss
                                                                     if (g[u][v]) continue;
int tsn = 0, tnn = 0;
                                                                     copy_n(all[d], an, all[d + 1]);
        else --t;
                                                                     all[d + 1][an] = v;
     }
                                                                     for (int j = 0; j < sn; ++j)</pre>
   }
                                                                       if (g[v][some[d][j]])
 }
                                                                         some[d + 1][tsn++] = some[d][j];
};
                                                                     for (int j = 0; j < nn; ++j)</pre>
      Maximum Clique [1ad4b2]
                                                                       if (g[v][none[d][j]])
                                                                         none[d + 1][tnn++] = none[d][j];
struct MaxClique { // fast when N <= 100</pre>
                                                                     dfs(d + 1, an + 1, tsn, tnn);
  bitset<N> G[N], cs[N];
                                                                     some[d][i] = 0, none[d][nn++] = v;
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
                                                                int solve() {
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
                                                                  iota(some[0], some[0] + n, 1);
                                                                   S = 0, dfs(0, 0, n, 0);
  void add_edge(int u, int v) {
                                                                   return S;
   G[u][v] = G[v][u] = 1;
                                                              };
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
                                                              5.10 Minimum Mean Cycle [3e5d2b]
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(iter(r), [\&](int x, int y) \{ return d[x] > d \}
                                                               // from 8BQube
          [y]; });
                                                              11 road[N][N]; // input here
                                                              struct MinimumMeanCycle {
    }
                                                                11 dp[N + 5][N], n;
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
                                                                pll solve() {
    cs[1].reset(), cs[2].reset();
                                                                   ll a = -1, b = -1, L = n + 1;
                                                                   for (int i = 2; i <= L; ++i)</pre>
    for (int p : r) {
                                                                     for (int k = 0; k < n; ++k)
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
                                                                       for (int j = 0; j < n; ++j)
      if (k > rgt) cs[++rgt + 1].reset();
                                                                         dp[i][j] =
      cs[k][p] = 1;
                                                                           min(dp[i - 1][k] + road[k][j], dp[i][j]);
                                                                   for (int i = 0; i < n; ++i) {</pre>
      if (k < lft) r[tp++] = p;
                                                                     if (dp[L][i] >= INF) continue;
                                                                     11 ta = 0, tb = 1;
for (int j = 1; j < n; ++j)</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
          ]._Find_next(p))
                                                                       if (dp[j][i] < INF &&</pre>
                                                                         {\sf ta} \ * \ ({\sf L} \ - \ {\sf j}) \ < \ ({\sf dp[L][i]} \ - \ {\sf dp[j][i]}) \ * \ {\sf tb})
        r[tp] = p, c[tp] = k, ++tp;
                                                                         ta = dp[L][i] - dp[j][i], tb = L - j;
    dfs(r, c, l + 1, mask);
                                                                     if (ta == 0) continue;
```

if (a == -1 || a \* tb > ta \* b) a = ta, b = tb;

**if** (a != -1) {

\_\_gcd(a, b);

return pll(a / g, b / g);

11 g =

```
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    return pll(-1LL, -1LL);
  void init(int _n) {
   n = _n;
for (int i = 0; i < n; ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
 }
};
       Minimum Steiner Tree [21acea]
5.11
// from 8BQube
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcost[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
      for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
   }
  void add_edge(int ui, int vi, int wi) {
    dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
          dst[i][j] =
            min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)</pre>
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
              submsk = (submsk - 1) \& msk)
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
              vcost[i]);
      for (int i = 0; i < n; ++i) {</pre>
        tdst[i] = INF;
        for (int j = 0; j < n; ++j)
          tdst[i] =
            min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    }
    int ans = INF;
    for (int i = 0; i < n; ++i)</pre>
     ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
 }
};
5.12 Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
  for (int y : D[x]) vis[y] = 1;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
  for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
  for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) c4 += vis[z]++;
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M))
```

```
Math
6.1 Extended Euclidean Algorithm [c51ae9]
// ax+ny = 1, ax+ny == ax == 1 \ (mod n)
void extgcd(ll x, ll y, ll &g, ll &a, ll &b) {
  if (y == 0) g = x, a = 1, b = 0;
  else extgcd(y, x \% y, g, b, a), b -= (x / y) * a;
6.2 Floor & Ceil [134881]
11 ifloor(ll a, ll b){
  return a / b - (a % b && (a < 0) ^ (b < 0));
ll iceil(ll a, ll b){
  return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre [4e4b23]
// the Jacobi symbol is a generalization of the
    Leaendre symbol
// such that the bottom doesn't need to be prime.
// (n|p) -> same as Legendre
// (n|ab) = (n|a)(n|b)
// work with Long Long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a \%= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  return s:
}
// 0: a == 0
// -1: a isn't a quad res of p
// else: return X with X^2 \% p == a
// doesn't work with long long
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
  int b, d;
  for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p
    )) % p;
f1 = (2LL * f0 * f1) % p;
```

# 6.4 Simplex [aa7741]

f0 = tmp;

return g0;

}

```
// maximize c^T x
// subject to Ax <= b, x >= 0
// and stores the solution;
typedef long double T; // Long double, Rational, double
     + mod<P>..
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-9, inf = 1/.0;
#define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
    #define rep(i, 1, n) for(int i = 1; i < n; i++)
```

b = 1LL \* b \* x % m;

```
for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
struct LPSolver {
  int m, n;
  vector<int> N, B;
                                                                     if (p.find(s) != p.end()) return i + kStep - p[s];
  vvd D:
                                                                   return -1;
  LPSolver(const vvd& A, const vd& b, const vd& c) :
    m(SZ(b)), n(SZ(c)), N(n+1), B(m), D(m+2, vd(n+2)) {
                                                                 int DiscreteLog(int x, int y, int m) {
       rep(i,0,m) rep(j,0,n) D[i][j] = A[i][j];
                                                                   if (m == 1) return 0;
       rep(i,0,m) \{ B[i] = n+i; D[i][n] = -1; D[i][n+1] \}
                                                                   int s = 1;
                                                                   for (int i = 0; i < 100; ++i) {
           = b[i];
                                                                     if (s == y) return i;
s = 1LL * s * x % m;
       rep(j,0,n) \{ N[j] = j; D[m][j] = -c[j]; \}
       N[n] = -1; D[m+1][n] = 1;
                                                                   if (s == y) return 100;
  void pivot(int r, int s) {
                                                                   int p = 100 + DiscreteLog(s, x, y, m);
    T *a = D[r].data(), inv = 1 / a[s];
rep(i,0,m+2) if (i != r && abs(D[i][s]) > eps) {
                                                                   if (fpow(x, p, m) != y) return -1;
                                                                   return p; //returns: x^p = y \pmod{m}
       T *b = D[i].data(), inv2 = b[s] * inv;
       rep(j,0,n+2) b[j] -= a[j] * inv2;
                                                                 6.7 Miller Rabin & Pollard Rho [d3ecd2]
       b[s] = a[s] * inv2;
                                                                 // n < 4,759,123,141
                                                                                              3 : 2, 7, 61
    rep(j,0,n+2) if (j != s) D[r][j] *= inv;
                                                                 // n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
    rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
                                                                 // n < 3,474,749,660,383 6 : primes <= 13
    D[r][s] = inv;
                                                                 // n < 2^64
    swap(B[r], N[s]);
                                                                 // 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                 11 mul(l1 a, l1 b, l1 n){
                                                                   return (__int128)a * b % n;
  bool simplex(int phase) {
    int x = m + phase - 1;
                                                                 bool Miller_Rabin(ll a, ll n) {
    for (;;) {
                                                                   if ((a = a % n) == 0) return 1;
       int s = -1;
                                                                   if (n % 2 == 0) return n == 2;
       rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
                                                                   if (D[x][s] >= -eps) return true;
       int r = -1;
                                                                   for (; tmp; tmp >>= 1, a = mul(a, a, n))
       rep(i,0,m) {
                                                                     if (tmp & 1) x = mul(x, a, n);
         if (D[i][s] <= eps) continue;</pre>
                                                                   if (x == 1 || x == n - 1) return 1;
         if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])
                                                                   while (--t)
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
                                                                     if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                   return 0;
       if (r == -1) return false;
       pivot(r, s);
                                                                 bool prime(ll n){
    }
                                                                   vector<11> tmp = {2, 325, 9375, 28178, 450775,}
  }
                                                                        9780504, 1795265022};
                                                                   for(ll i : tmp)
  T solve(vd &x) {
                                                                     if(!Miller_Rabin(i, n)) return false;
    int r = 0;
                                                                   return true;
    rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] < -eps) {</pre>
                                                                 map<ll, int> cnt;
       pivot(r, n);
                                                                 void PollardRho(ll n) {
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                   if (n == 1) return;
                                                                   if (prime(n)) return ++cnt[n], void();
       rep(i,0,m) if (B[i] == -1) {
                                                                   if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
                                                                        void();
         rep(j,1,n+1) ltj(D[i]);
                                                                   11 x = 2, y = 2, d = 1, p = 1;
         pivot(i, s);
                                                                 #define f(x, n, p) ((mul(x, x, n) + p) % n)
      }
                                                                   while (true) {
                                                                     if (d != n && d != 1) {
    bool ok = simplex(1); x = vd(n);
                                                                       PollardRho(n / d);
    rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
                                                                       PollardRho(d);
    return ok ? D[m][n+1] : inf;
                                                                       return;
                                                                     if (d == n) ++p;
                                                                     x = f(x, n, p), y = f(f(y, n, p), n, p);
d = gcd(abs(x - y), n);
      Simplex Construction
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for
all 1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
                                                                 }
1. In case of minimization, let c_i' = -c_i
                                                                 6.8 XOR Basis [006505]
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n} A_{ji} x_i = b_j \rightarrow \mathsf{add} \subseteq \mathsf{and} \supseteq.
                                                                 const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                 struct Basis{
6.6 DiscreteLog [da27bf]
                                                                   int total = 0, rank = 0;
                                                                   vector<ll> b;
int DiscreteLog(int s, int x, int y, int m) {
                                                                   Basis(): b(digit) {}
  constexpr int kStep = 32000;
                                                                   bool add(ll v){ // Gauss Jordan Elimination
  unordered_map<int, int> p;
                                                                     total++;
  int b = 1;
                                                                     for(int i = digit - 1; i >= 0; i--){
  for (int i = 0; i < kStep; ++i) {</pre>
                                                                       if(!(1LL << i & v)) continue;</pre>
    p[y] = i;
                                                                       if(b[i] != 0){
    y = 1LL * y * x % m;
                                                                          v ^= b[i];
```

continue;

for (int i = 0; i < M; i++) {</pre>

if (ispiv[i]) continue;

```
vector<ll> h(M);
       for(int j = 0; j < i; j++)
  if(1LL << j & v) v ^= b[j];</pre>
                                                                                    h[i] = 1;
                                                                                    for (int j = 0; j < rk; j++)</pre>
        for(int j = i + 1; j < digit; j++)</pre>
                                                                                       h[piv[j]] = rref[j][i] ? MOD - rref[j][i] : 0;
          if(1LL << i & b[j]) b[j] ^= v;</pre>
                                                                                    homo.pb(h);
       b[i] = v;
       rank++;
                                                                               }
       return true;
                                                                            };
                                                                            6.10
                                                                                      Chinese Remainder Theorem [6ef4a3]
     return false;
                                                                            pll solve_crt(ll x1, ll m1, ll x2, ll m2){
  11 \text{ getmax}(11 \text{ x} = 0)
                                                                               ll g = gcd(m1, m2);
     for(ll i : b) x = max(x, x ^ i);
                                                                               if ((x2 - x1) % g) return {0, 0}; // no sol
     return x;
                                                                               m1 /= g; m2 /= g;
                                                                               11 _, p, q;
  11 \text{ getmin}(11 \text{ x} = 0){
                                                                               extgcd(m1, m2, _, p, q); // p <= C

ll lcm = m1 * m2 * g;

ll res = ((__int128)p * (x2 - x1) % lcm * m1 % lcm +
     for(11 i : b) x = min(x, x ^ i);
     return x;
                                                                                    x1) % lcm;
  bool can(ll x){
                                                                               // be careful with overflow, C^3
     return getmin(x) == 0;
                                                                               return {(res + lcm) % lcm, lcm}; // (x, m)
                                                                            }
  11 kth(ll k){ // kth smallest, 0-indexed
     vector<11> tmp;
                                                                            6.11 Sqrt Decomposition [8d7bc0]
     for(ll i : b) if(i) tmp.pb(i);
     11 \text{ ans} = 0;
                                                                            // for all i in [l, r], floor(n / i) = x
     for(int i = 0; i < SZ(tmp); i++)</pre>
                                                                            for(int l = 1, r; l <= n; l = r + 1){
       if(1LL << i & k) ans ^= tmp[i];</pre>
                                                                               int x = ifloor(n, 1);
                                                                               r = ifloor(n, x);
     return ans:
};
                                                                            // for all i in [l, r], ceil(n / i) = x
                                                                            for(int 1, r = n; r >= 1; r = 1 - 1){
6.9 Linear Equation [056191]
                                                                               int x = iceil(n, r);
vector<int> RREF(vector<vector<ll>> &mat) { // SCOPE
                                                                               l = iceil(n, x);
  int N = SZ(mat), M = SZ(mat[0]);
                                                                            6.12 Floor Sum
  int rk = 0;
                                                                             • m = \lfloor \frac{an+b}{2} \rfloor
  vector<int> cols;
  for (int i = 0; i < M; i++) {</pre>
                                                                             • Time complexity: O(\log n)
     int cnt = -1;
     for (int j = N - 1; j >= rk; j--)
                                                                                 f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
       if(mat[j][i] != 0) cnt = j;
     if (cnt == -1) continue;
                                                                                                \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), \end{cases}
     swap(mat[rk], mat[cnt]);
     11 lead = mat[rk][i];
                                                                                                                                  a \ge c \lor b \ge c
     for (int j = 0; j < M; j++) mat[rk][j] = mat[rk][j]</pre>
                                                                                                                                  n < 0 \lor a = 0
            * inv(lead) % MOD;
                                                                                                 nm - f(c, c - b - 1, a, m - 1), otherwise
     for (int j = 0; j < N; j++) {</pre>
       if (j == rk) continue;
                                                                            g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
       11 tmp = mat[j][i];
       for (int k = 0; k < M; k++)
          mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % MOD
                                                                                            \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2}
                 + MOD) % MOD;
                                                                                            +g(a \bmod c, b \bmod c, c, n),
                                                                                                                                        a \ge c \lor b \ge c
                                                                                                                                        n < 0 \lor a = 0
     cols.pb(i);
                                                                                            \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
     rk++;
  }
                                                                                             -h(c, c-b-1, a, m-1)),
                                                                                                                                        otherwise
  return cols;
                                                                            h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2
struct LinearEquation { // SCOPE HASH
  bool ok:
                                                                                            \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
  vector<11> par; //particular solution (Ax = b)
                                                                                            +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot \tilde{n}(n+1)
  vector<vector<ll>> homo; //homogenous (Ax = 0)
  vector<vector<ll>> rref;
                                                                                            +h(a \bmod c, b \bmod c, c, n)
  //first M columns are matrix A
                                                                                            +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
  //last column of eq is vector b
                                                                                            +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                                                        a \ge c \lor b \ge c
  void solve(const vector<vector<11>>> &eq) {
                                                                                            0.
                                                                                                                                        n < 0 \lor a = 0
     int M = SZ(eq[0]) - 1;
     rref = eq;
                                                                                            nm(m+1) - 2g(c, c-b-1, a, m-1)
     auto piv = RREF(rref);
                                                                                            -2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), otherwise
     int rk = piv.size();
     if(piv.size() && piv.back() == M)
                                                                                  Polynomial
                                                                            7
       return ok = 0, void();
     ok = 1;
                                                                            7.1 FWHT [c9cdb6]
     par.resize(M);
                                                                            /* x: a[j], y: a[j + (L >> 1)]
     vector<bool> ispiv(M);
     for (int i = 0;i < rk;i++) {</pre>
                                                                            or: (y += x * op), and: (x += y * op)
                                                                            xor: (x, y = (x + y) * op, (x - y) * op)
       par[piv[i]] = rref[i][M];
                                                                            invop: or, and, xor = -1, -1, 1/2 */
       ispiv[piv[i]] = 1;
```

void fwt(int \*a, int n, int op) { //or

for (int L = 2; L <= n; L <<= 1)
 for (int i = 0; i < n; i += L)</pre>

```
for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
    ];
void subset_convolution(int *a, int *b, int *c, int L)
  // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)</pre>
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i) fwt(h[i], n, -1);</pre>
                                                                     }
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
7.2 FFT [13ec2f]
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double
                                                                 }
using val_t = complex<double>;
                                                               };
template<int MAXN>
struct FFT {
  const double PI = acos(-1);
  val_t w[MAXN];
                                                                    ++i)
    for (int i = 0; i < MAXN; ++i) {
  double arg = 2 * PI * i / MAXN;</pre>
      w[i] = val_t(cos(arg), sin(arg));
    }
  void bitrev(vector<val_t> &a, int n) //same as NTT
  void trans(vector<val_t> &a, int n, bool inv = false)
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
           val_t + mp = a[j + dl] * (inv ? conj(w[x]) : w
               [x]);
           a[i + d1] = a[j] - tmp;
          a[j] += tmp;
        }
      }
    if (inv) {
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
  //multiplying two polynomials A * B:
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
  //A[i] *= B[i], fft.trans(A, siz, 1);
7.3 NTT [bf683f]
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
// only works when sz(A) + sz(B) - 1 <= MAXN
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
struct NTT {
  11 w[MAXN];
  11 mpow(11 a, 11 n);
  11 minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
  }
```

```
void bitrev(vector<ll> &a, int n) {
    int i = 0:
    for (int j = 1; j < n - 1; ++j) {</pre>
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
  void operator()(vector<ll> &a, int n, bool inv =
      false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
            ) {
          11 tmp = a[j + d1] * w[x] % P;
          if ((a[j + dl] = a[j] - tmp) < 0) a[j + dl]
          if ((a[j] += tmp) >= P) a[j] -= P;
        }
    if (inv) {
      reverse(a.begin()+1, a.begin()+n);
      11 invn = minv(n);
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
7.4 Polynomial Operation [77a8a8]
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
#define neg(x) (x ? P - x : 0)
#define V (*this)
template <int MAXN, 11 P, 11 RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<ll>::vector;
  static inline NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
      .data(), min(p.n(), m), data()); }
  Poly &irev() { return reverse(data(), data() + n()),
      V; }
  Poly &isz(int m) { return resize(m), V; }
  static ll minv(ll x) { return ntt.minv(x); }
  // == PART HASH ==
  Poly &iadd(const Poly &rhs) { // SCOPE HASH
    fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
    return V; // need n() == rhs.n()
  Poly &imul(11 k) { // SCOPE HASH
    fi(0, n()) V[i] = V[i] * k % P;
    return V;
  Poly Mul(const Poly &rhs) const { // SCOPE HASH
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    assert(m <= MAXN);</pre>
    Poly X(V, m), Y(rhs, m);
    ntt(X, m), ntt(Y, m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X, m, true);
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // SCOPE HASH
    if (n() == 1) return {minv(V[0])};
    int m = 1; // need V[0] != 0, 2*sz<=MAXN</pre>
    while (m < n() * 2) m <<= 1;</pre>
    assert(m <= MAXN);
    Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);
    Poly Y(V, m);
    ntt(Xi, m), ntt(Y, m);
    fi(0, m) {
      Xi[i] *= (2 - Xi[i] * Y[i]) % P;
      if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    ntt(Xi, m, true);
    return Xi.isz(n());
```

```
Poly &shift_inplace(const 11 &c) { // SCOPE HASH
  int n = V.n(); // 2 * sz <= MAXN</pre>
  vector<ll> fc(n), ifc(n);
  fc[0] = ifc[0] = 1;
  for (int i = 1; i < n; i++) {</pre>
    fc[i] = fc[i - 1] * i % P;
    ifc[i] = minv(fc[i]);
  for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
  Poly g(n);
  11 cp = 1;
  for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
       cp = cp * c % P;
  V = V.irev().Mul(g).isz(n).irev();
  for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
  return V;
}
// == PART HASH ==
Poly shift(const 11 &c) const { return Poly(V).
shift_inplace(c); }
Poly _Sqrt() const { // Jacobi(V[0], P) = 1
if (n() == 1) return {QuadraticResidue(V[0], P)};
  Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
      1);
// == PART HASH ==
Poly Sqrt() const { // SCOPE HASH
  Poly a; // 2 * sz <= MAXN
  bool has = 0;
  for (int i = 0; i < n(); i++) {</pre>
    if (V[i]) has = 1;
    if (has) a.push_back(V[i]);
  if (!has) return V;
  if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
    return Poly();
  a = a.isz((n() + a.n()) / 2)._Sqrt();
  int sz = a.n();
  a.isz(n());
  rotate(a.begin(), a.begin() + sz, a.end());
  return a:
pair<Poly, Poly> DivMod(const Poly &rhs) const { //
    SCOPE HASH
  if (n() < rhs.n()) return {{0}, V};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs); // (rhs.)back() != 0
  X.irev().isz(m);
  Poly Y(V);
  Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
  X = rhs.Mul(Q), Y = V;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
// == PART HASH ==
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
  return ret:
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
// == PART HASH ==
vector<ll> _eval(const vector<ll> &x, const vector<
    Poly> &up) const { // SCOPE HASH
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
```

```
// fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
      .second:
  down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
      _tmul(m, V);
  fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
      1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
static vector<Poly> _tree1(const vector<ll> &x) { //
    SCOPE HASH
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = {neg(x[i]), 1};
  for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
      Mul(up[i * 2 + 1]);
  return up:
vector<ll> Eval(const vector<ll> &x) const { // 1e5,
    15
  auto up = _tree1(x);
  return _eval(x, up);
static Poly Interpolate(const vector<11> &x, const
    vector<ll> &y) { // SCOPE HASH
  const int m = (int)x.size(); // 1e5, 1.4s
  vector<Poly> up = _tree1(x), down(m * 2);
  vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
 for (int i = m - 1; i > 0; --i)
  down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
       down[i * 2 + 1].Mul(up[i * 2]));
 return down[1];
}
// == PART HASH ==
Poly Ln() const { // V[0] == 1, 2*sz<=MAXN
 return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { //V[0] == 0,2*sz <= MAXN
  if (n() == 1) return {1};
  Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln();
  Y[0] = P - 1;
  fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
  return X.Mul(Y).isz(n());
// == PART HASH ==
// M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(11 k) const { // 2*sz<=MAXN // SCOPE HASH
  int nz = 0;
  while (nz < n() && !V[nz]) ++nz;</pre>
  if (nz * min(k, (11)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly{1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const ll c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
      n()).irev();
// sum_j w_j [x^j] f(x^i) for i \in [0, m]
Poly power_projection(Poly wt, int m) { // SCOPE HASH
  assert(n() == wt.n()); // 4*sz <= MAXN!
  if (!n()) {
    return Poly(m + 1, 0);
  if (V[0] != 0) {
    11 c = V[0];
    V[0] = 0;
    Poly A = V.power_projection(wt, m);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
       factorial
    Poly B(m + 1);
    11 pow = 1;
    fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
        * c % P; // inv. of fac
    A = A.Mul(B).isz(m + 1);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P;
    return A;
  int n = 1;
```

```
while (n < V.n()) n *= 2;</pre>
  isz(n), wt.isz(n).irev();
  int k = 1;
  Poly p(wt, 2 * n), q(V, 2 * n);
  q.imul(P - 1);
  while (n > 1) {
    Poly r(2 * n * k);
    fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
        q[i]));
    Poly pq = p.Mul(r).isz(4 * n * k);
    Poly qq = q.Mul(r).isz(4 * n * k);
    fi(0, 2 * n * k) {
      pq[2 * n * k + i] += p[i];
      qq[2 * n * k + i] += q[i] + r[i];
      pq[2 * n * k + i] %= P;
      qq[2 * n * k + i] \% = P;
    fill(p.begin(), p.end(), 0);
    fill(q.begin(), q.end(), 0);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
   p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];
   q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];
    n /= 2, k *= 2;
  Poly ans(k);
  fi(0, k) ans[i] = p[2 * i];
  return ans.irev().isz(m + 1);
Poly FPSinv() { // SCOPE HASH
  const int n = V.n() - 1;
  if (n == -1) return {};
  assert(V[0] == 0);
  if (n == 0) return V;
  assert(V[1] != 0);
  ll c = V[1], ic = minv(c);
  imul(ic);
  Poly wt(n + 1);
  wt[n] = 1;
  Poly A = V.power_projection(wt, n);
  Poly g(n);
  fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
  g = g.Pow(neg(minv(n)));
  g.insert(g.begin(), 0);
  11 pow = 1;
  fi(0, g.n()) g[i] = g[i] * pow % P, pow = pow * ic
      % P;
  return g;
Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
    ] = k; // SCOPE HASH
  return Poly(*this).irev().Mul(rhs).isz(n()).irev();
Poly FPScomp(Poly g) { // solves V(g(x)) // SCOPE
  auto rec = [&](auto &rec, int n, int k, Poly Q) ->
      Poly {
    if (n == 1) {
      Poly p(2 * k);
      irev();
      fi(0, k) p[2 * i] = V[i];
      return p;
    Poly R(2 * n * k);
    fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
        Q[i]));
    Poly QQ = Q.Mul(R).isz(4 * n * k);
fi(0, 2 * n * k) {
      QQ[2 * n * k + i] += Q[i] + R[i];
      QQ[2 * n * k + i] %= P;
    Poly nxt_Q(2 * n * k);
for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
      nxt_{Q[n * j + i]} = QQ[(2 * n) * j + (2 * i + 0)
           ];
    Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
Poly pq(4 * n * k);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
```

```
pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +
        i];
pq[(2 * n) * j + (2 * i + 1)] %= P;
      Poly p(2 * n * k);
      fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i
          ]) % P;
      pq.pop_back();
      Poly x = pq.TMul(R);
      fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
      return p;
    };
    int sz = 1:
    while(sz < n() || sz < g.n()) sz <<= 1;
    return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
         sz)).isz(sz).irev();
  }
}:
#undef fi
#undef V
#undef neg
using Poly_t = Poly<1 << 19, 998244353, 3>;
```

# 7.5 Generating Function Ordinary Generating Function

- C(x) = A(rx):  $c_n = r^n a_n$  的一般生成函數。
- C(x) = A(x) + B(x):  $c_n = a_n + b_n$  的一般生成函數。
- C(x) = A(x)B(x):  $c_n = \sum_{i=0}^n a_i b_{n-i}$  的一般生成函數。  $C(x) = A(x)^k$ :  $c_n = \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$  的一般生成函數。
- C(x) = xA(x)':  $c_n = na_n$  的一般生成函數。
- $C(x) = \frac{A(x)}{1-x}$ :  $c_n = \sum_{i=0}^n a_i$  的一般生成函數。
- $C(x) = A(1) + x \frac{A(1) A(x)}{1 x}$ :  $c_n = \sum_{i=1}^{\infty} a_i$  的一般生成函數。

### 常用展開式

- $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$ ,  ${a \choose n} = \frac{a(a-1)(a-2)\dots(a-n+1)}{n!}$ .

• 卡特蘭數  $f(x) = \frac{1 - \sqrt{1 - 4x}}{2\pi}$ 

# **Exponential Generating Function**

 $a_0, a_1, \ldots$  的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$ :  $c_n = a_n + b_n$  的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$ :  $c_n = a_{n+k}$  的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$ :  $c_n = na_n$  的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$ :  $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$  的指數生成函數
- $\sum_{i_1+i_2+\cdots+i_k=n}^{\sum_{k=0}^{n}\binom{i_1}{n}}\binom{n}{i_1,i_2,\ldots,i_k}a_ia_{i_2}\ldots a_{i_k}$  的指數生成函數 •  $\hat{C}(x) = \hat{A}(x)^k$ :
- $\hat{C}(x) = \exp(A(x))$ : 假設 A(x) 是一個分量 (component) 的生成函數,那  $\hat{C}(x)$  是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 F(0),G(0)=0,  $F'(0),G'(0)\neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$
$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

# 7.6 Bostan Mori [41c3bc]

```
const 11 mod = 998244353;
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = \{P, Q\};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
```

/\* center i: radius z[i \* 2 + 1] / 2

center i, i + 1: radius z[i \* 2 + 2] / 2

```
vector<11> Qn(sz);
                                                                 both aba, abba have radius 2 */
  vector N(2, vector<ll>(sz));
                                                              vector<int> manacher(const string &tmp){ // 0-based
                                                                string s = "%";
  while(k) {
    fill(iter(Qn), 0);
for(int i = 0; i < d + 1; i++){
                                                                int 1 = 0, r = 0;
                                                                for(char c : tmp) s += c, s += '%';
      Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);
                                                                vector<int> z(SZ(s));
                                                                for(int i = 0; i < SZ(s); i++){
  z[i] = r > i ? min(z[2 * l - i], r - i) : 1;
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    }
                                                                  while(i - z[i] >= 0 \&\& i + z[i] < SZ(s)
    ntt(Qn, sz, false);
                                                                         && s[i + z[i]] == s[i - z[i]])
                                                                    ++z[i];
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){
                                                                  if(z[i] + i > r) r = z[i] + i, l = i;
      fill(iter(N[i]), 0);
                                                                }
      copy(iter(M[i]), N[i].begin());
                                                                return z;
                                                             }
      ntt(N[i], sz, false);
      for(int j = 0; j < sz; j++)</pre>
                                                             8.3
                                                                    Lyndon Factorization [7c612b]
        N[i][j] = N[i][j] * Qn[j] % mod;
      ntt(N[i], sz, true);
                                                              // partition s = w[0] + w[1] + ... + w[k-1],
      for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
                                                              // w[0] >= w[1] >= ... >= w[k-1]
        M[i][j >> 1] = N[i][j];
                                                              // each w[i] strictly smaller than all its suffix
                                                              void duval(const string &s, vector<pii> &w) {
                                                                for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
    k >>= 1;
                                                                  for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
  }
                                                                    k = (s[k] < s[j] ? i : k + 1);
  return M[0][0] * ntt.minv(M[1][0]) % mod;
                                                                  // if (i < n / 2 && j >= n / 2) {
                                                                  // for min cyclic shift, call duval(s + s)
11 LinearRecursion(vector<ll> a, vector<ll> c, ll k) {
                                                                  // then here s.substr(i, n / 2) is min cyclic shift
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
                                                                  for (; i <= k; i += j - k)</pre>
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
                                                                    w.pb(pii(i, j - k)); // s.substr(l, len)
      1)));
                                                             }
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
                                                             8.4 Suffix Array [cd67ea]
  auto A = a; A.resize(sz);
                                                              struct SuffixArray {
  auto C = c; C.resize(sz);
                                                                vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]
  ntt(A, sz, false), ntt(C, sz, false);
                                                                    and sa[i-1]
  for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
                                                                                             // sa[0] = s.size()
  ntt(A, sz, true);
                                                                                             // character should be 1-
  A.resize(d);
                                                                                                 hased
                                                                SuffixArray(string& s, int lim=256) { // or
  return BostanMori(A, c, k);
                                                                    basic_string<int>
                                                                  int n = s.size() + 1, k = 0, a, b;
                                                                  vector<int> x(n, 0), y(n), ws(max(n, lim));
     String
8
                                                                  rank.assign(n, 0);
                                                                  for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
      KMP Algorithm [c8b75f]
                                                                  sa = lcp = y, iota(sa.begin(), sa.end(), 0);
                                                                  for (int j = 0, p = 0; p < n; j = max(1, j * 2),
// 0-based
                                                                      lim = p) {
// fail[i] = max k<i s.t. s[0..k] = s[i-k..i]
                                                                    p = j, iota(y.begin(), y.end(), n - j);
vector<int> kmp_build_fail(const string &s){
                                                                    for (int i = 0; i < n; i++)</pre>
  int n = SZ(s);
                                                                      if (sa[i] >= j) y[p++] = sa[i] - j;
  vector<int> fail(n, -1);
                                                                    for (int &i : ws) i = 0;
  int cur = -1;
                                                                    for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
  for(int i = 1; i < n; i++){</pre>
                                                                    for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];</pre>
    while(cur != -1 && s[cur + 1] != s[i])
      cur = fail[cur];
                                                                    swap(x, y), p = 1, x[sa[0]] = 0;
    if(s[cur + 1] == s[i])
                                                                    for(int i = 1; i < n; i++){</pre>
      cur++;
                                                                      a = sa[i - 1], b = sa[i];
    fail[i] = cur;
                                                                      x[b] = (y[a] == y[b] && y[a + j] == y[b + j])?
                                                                            p - 1 : p++;
  return fail;
                                                                  for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
void kmp_match(const string &s, const vector<int> &fail
                                                                  for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
    , const string &t){
                                                                    for (k \&\& k--, j = sa[rank[i] - 1];
  int cur = -1;
                                                                        s[i + k] == s[j + k]; k++);
  int n = SZ(s), m = SZ(t);
                                                                }
  for(int i = 0; i < m; i++){</pre>
                                                             };
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t
        [i]))
                                                              8.5 Suffix Automaton [016373]
      cur = fail[cur];
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
                                                              struct exSAM {
                                                                const int CNUM = 26;
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
                                                                // len: maxlength, link: fail link
  }
                                                                // LenSorted: topo order, cnt: occur
                                                                vector<int> len, link, lenSorted, cnt;
8.2 Manacher Algorithm [caf0f4]
                                                                vector<vector<int>> next;
                                                                int total = 0;
```

int newnode() { return total++;

const int n = s.size();

if (n == 1) return;

```
void init(int n) { // total number of characters
    len.assign(2 * n, 0); link.assign(2 * n, 0);
    lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
next.assign(2 * n, vector<int>(CNUM));
    newnode(), link[0] = -1;
  int insertSAM(int last, int c) {
    // not exSAM: cur = newnode(), p = Last
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len[p] + 1 == len[q]) return link[cur] = q, cur
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[clone][i] = len[next[q][i]] ? next[q][i] :
          0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  }
  void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
    }
  }
  void build() {
    queue<int> q;
                                                             }
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    vector<int> lc(total);
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
        i]]] = i;
  }
  void solve() {
    for (int i = total - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
8.6 Z-value Algorithm [488d87]
// z[i] = max \ k \ s.t. \ s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z_function(const string &s){
  int n = s.size();
  vector<int> z(n);
  for(int i = 1, l = 0, r = 0; i < n; i++){</pre>
    if(i <= r) z[i] = min(r - i + 1, z[i - 1]);</pre>
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
      z[i]++;
    if(i + z[i] - 1 > r)
      \hat{l} = i, r = i + z[i] - 1;
  return z;
8.7 Main Lorentz [fcfb8f]
struct Rep{ int minl, maxl, len; };
vector<Rep> rep; // 0-base
```

 $// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)$ void main\_lorentz(const string &s, int sft = 0) {

```
const int nu = n / 2, nv = n - nu;
  const string u = s.substr(0, nu), v = s.substr(nu),
         ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
             ());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const auto z1 = z_function(ru), z2 = z_function(v +
              z3 = z_function(ru + '#' + rv), z4 =
                  z_function(v);
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
  auto add_rep = [&](bool left, int c, int l, int k1,
      int k2) {
    const int L = max(1, 1 - k2), R = min(1 - left, k1)
    if (L > R) return;
    if (left) rep.emplace_back(Rep({sft + c - R, sft +
         c - L, 1}));
    else rep.emplace_back(Rep({sft + c - R - l + 1, sft
          + c - L - 1 + 1, 1));
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int 1, k1, k2;
    if (cntr < nu) {</pre>
      1 = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
8.8
     AC Automaton [f529e6]
const int SIGMA = 26;
struct AC_Automaton {
  // child: trie, next: automaton
  vector<vector<int>> child, next;
  vector<int> fail, cnt, ord;
  int total = 0;
  int newnode() {
    return total++;
  void init(int len) { // len >= 1 + total len
    child.assign(len, vector<int>(26, -1));
next.assign(len, vector<int>(26, -1));
    fail.assign(len, -1); cnt.assign(len, 0);
    ord.clear();
    newnode();
  int input(string &s) {
    int cur = 0;
    for (char c : s) {
      if (child[cur][c - 'A'] == -1)
    child[cur][c - 'A'] = newnode();
cur = child[cur][c - 'A'];
    return cur; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(0), fail[0] = -1;
    while(!q.empty()) {
      int R = q.front();
      q.pop(); ord.pb(R);
      for (int i = 0; i < SIGMA; i++)</pre>
        if (child[R][i] != -1) {
           int X = next[R][i] = child[R][i], Z = fail[R
           while (Z != -1 && child[Z][i] == -1)
            Z = fail[Z];
           fail[X] = Z != -1 ? child[Z][i] : 0;
           q.push(X);
        else next[R][i] = R ? next[fail[R]][i] : 0;
```

```
}
 }
  void solve() {
    for (int i : ord | views::reverse)
      if (i) cnt[fail[i]] += cnt[i];
};
```

#### Palindrome Automaton [8a071b] 8.9

```
struct PalindromicTree {
  struct node {
    int nxt[26], fail, len; // num = depth of fail link
    int cnt, num; // cnt = occur, num = #pal_suffix of
        this node
    node(int 1 = 0) : nxt{}, fail(0), len(1), cnt(0), num
        (0) \{ \}
 };
  vector<node> st; vector<int> s; int last, n;
  void init() {
    st.clear(); s.clear(); last = 1; n = 0;
    st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
  int getFail(int x) {
    while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x:
  void add(int c) {
    s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
      int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
      st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
  void dpcnt() {
    for(int i = SZ(st) - 1; i >= 0; i--){
      auto nd = st[i];
      st[nd.fail].cnt += nd.cnt;
  int size() { return (int)st.size() - 2; }
};
```

#### 9 Misc

# Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int 1 = 0, r = n; bool rv = pred(1, 0);
  while (r - l > 1) {
    int m = (1 + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else 1 = m;
  return pred(1, r % n) ? 1 : r % n;
```

# 9.2 Matroid

 $M=(E,\mathcal{I})$ , where  $\mathcal{I}\subseteq 2^E$  is nonempty, is a matroid if:

• If  $S \in \mathcal{I}$  and  $S' \subsetneq S$ , then  $S' \in \mathcal{I}$ .

• For  $S_1, S_2 \in \mathcal{I}$  s.t.  $|S_1| < |S_2|$ , there exists  $e \in S_2 \setminus S_1$  s.t.  $S_1 \cup \{e\} \in \mathcal{I}$ . Matroid intersection:

Start from  $S = \emptyset$ . In each iteration, let

•  $Y_1 = \{ x \notin S \mid S \cup \{x\} \in \mathcal{I}_1 \}$ 

•  $Y_2 = \{ x \notin S \mid S \cup \{x\} \in \mathcal{I}_2 \}$ 

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \notin S$ , create edges

•  $x \to y \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_1.$ 

•  $y \to x \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_2.$ 

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of  ${\cal S}$  will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x)if  $x \notin S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.

# 9.3 Simulate Annealing [ff826c]

```
ld anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<ld> rnd(0, 1);
  const ld dT = 0.001:
  // Argument p
  ld S_cur = calc(p), S_best = S_cur;
  for (ld T = 2000; T > eps; T -= dT) {
    // Modify p to p_prime
    const ld S_prime = calc(p_prime);
    const ld delta_c = S_prime - S_cur;
    ld prob = min((ld)1, exp(-delta_c / T));
    if (rnd(rnd_engine) <= prob)</pre>
      S_cur = S_prime, p = p_prime;
    if (S_prime < S_best) // find min</pre>
      S_best = S_prime, p_best = p_prime;
  return S_best;
```

# 9.4 Binary Search On Fraction [f6b9ec]

```
struct Q {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p * d, q + b.q * d};
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(11 N, auto &&pred) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step /= 2 : step *=</pre>
        2);)
      if (Q mid = hi.go(lo, len + step);
          mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

# 9.5 Min Plus Convolution [09b5c3]

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector
     <int> &b) {
  int n = SZ(a), m = SZ(b);
  vector<int> c(n + m - 1, INF);
auto dc = [&](auto Y, int l, int r, int jl, int jr) {
    if (1 > r) return;
    int mid = (1 + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)
  if (int i = mid - j; i >= 0 && i < n)</pre>
         if (best > a[i] + b[j])
           best = a[i] + b[j], from = j;
    Y(Y, l, mid - 1, jl, from), Y(Y, mid + 1, r, from,
         jr);
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
```

#### 10 Notes

### 10.1 Geometry **Rotation Matrix**

```
/cos θ
              -\sin\theta
\sin \theta
               \cos \theta
```

- rotate  $90^{\circ}$ :  $(x,y) \rightarrow (-y,x)$
- rotate  $-90^{\circ}$ :  $(x,y) \rightarrow (y,-x)$

# **Triangles**

Side lengths: a,b,c

Semiperimeter:  $p = \frac{a+b+c}{c}$ 

Area: 
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

Circumradius:  $R = \frac{abc}{4A}$ 

Inradius:  $r = \frac{A}{p}$ 

Length of median (divides triangle into two equal-area triangles):  $m_a =$ 

Length of bisector (divides angles in two): 
$$s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}$$

Law of sines: 
$$\frac{\sin\alpha}{a}=\frac{\sin\beta}{b}=\frac{\sin\gamma}{c}=\frac{1}{2R}$$
 Law of cosines: 
$$a^2=b^2+c^2-2bc\cos\alpha$$

Law of tosines. 
$$a = b + c = 2ac \cos \alpha$$

$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

# Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle  $\theta$ , area A and magic flux  $F = b^2 + d^2 - a^2 - c^2$ :

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is  $180^{\circ}$ , ef = ac + bd, and  $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$ .

## Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \operatorname{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \operatorname{atan2}(u, x) \end{array}$$

# Green's Theorem

$$\iint_{D} \left( \frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy = \oint_{L^{+}} (Pdx + Qdy)$$
 
$$\mathsf{Area} = \frac{1}{2} \oint_{I} x \, dy - y \, dx$$

Circular sector:

$$\begin{aligned} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r(r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{aligned}$$

# **Point-Line Duality**

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- $p_1, p_2, p_3$  are collinear  $\iff p_1^*, p_2^*, p_3^*$  intersect at a point p lies above  $l \iff l^*$  lies above  $p^*$
- lower convex hull ↔ upper envelope

# Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x})$$
 
$$\cosh x = \frac{1}{2}(e^x + e^{-x})$$
 
$$\cos n\pi = (-1)^n$$

$$\begin{split} \sin(\alpha+\beta) &= \sin\alpha\cos\beta + \cos\alpha\sin\beta\\ \cos(\alpha+\beta) &= \cos\alpha\cos\beta - \sin\alpha\sin\beta\\ \sin(2\alpha) &= 2\cos\alpha\sin\alpha\\ \cos(2\alpha) &= \cos^2\alpha - \sin^2\alpha\\ &= 2\cos^2\alpha - 1\\ &= 1 - 2\sin^2\alpha \end{split}$$

$$\begin{split} \tan(\alpha+\beta) &= \frac{\tan\alpha + \tan\beta}{1-\tan\alpha\tan\beta} \\ \sin\alpha + \sin\beta &= 2\sin\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \cos\alpha + \cos\beta &= 2\cos\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \sin\alpha\sin\beta &= \frac{1}{2}(\cos(\alpha-\beta)-\cos(\alpha+\beta)) \\ \sin\alpha\cos\beta &= \frac{1}{2}(\sin(\alpha+\beta)+\sin(\alpha-\beta)) \\ \cos\alpha\sin\beta &= \frac{1}{2}(\sin(\alpha+\beta)-\sin(\alpha-\beta)) \\ \cos\alpha\cos\beta &= \frac{1}{2}(\cos(\alpha-\beta)+\cos(\alpha+\beta)) \\ (V+W)\tan(\alpha-\beta)/2 &= (V-W)\tan(\alpha+\beta)/2 \end{split}$$

where V,W are lengths of sides opposite angles  $\alpha,\beta$ .

$$a\cos x + b\sin x = r\cos(x - \phi)$$
$$a\sin x + b\cos x = r\sin(x + \phi)$$

where 
$$r = \sqrt{a^2 + b^2}, \phi = \text{atan2}(b, a)$$
.

# 10.3 Calculus

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\begin{split} \frac{d}{dx} \arcsin x &= \frac{1}{\sqrt{1-x^2}} & \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^2}} \\ \frac{d}{dx} \tan x &= 1 + \tan^2 x & \frac{d}{dx} \arctan x = \frac{1}{1+x^2} \\ \int \tan ax &= -\frac{\ln|\cos ax|}{a} & \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2} \\ \int e^{-x^2} &= \frac{\sqrt{\pi}}{2} \mathrm{erf}(x) & \int x e^{ax} &= \frac{e^{ax}}{a^2} (ax-1) \\ \int \sin^2(x) &= \frac{x}{2} - \frac{1}{4} \sin 2x & \int \sin^3 x &= \frac{1}{12} \cos 3x - \frac{3}{4} \cos x \\ \int \cos^2(x) &= \frac{x}{2} + \frac{1}{4} \sin 2x & \int \cos^3 x &= \frac{1}{12} \sin 3x + \frac{3}{4} \sin x \\ \int x \sin x &= \sin x - x \cos x & \int x \cos x &= \cos x + x \sin x \\ \int x e^x &= e^x (x-1) & \int x^2 e^x &= e^x (x^2 - 2x + 2) \end{split}$$

$$\int x^2 \sin x = 2x \sin x - (x^2 - 2) \cos x$$

$$\int x^2 \cos x = 2x \cos x + (x^2 - 2) \sin x$$

$$\int e^x \sin x = \frac{1}{2} e^x (\sin x - \cos x)$$

$$\int e^x \cos x = \frac{1}{2} e^x (\sin x + \cos x)$$

$$\int x e^x \sin x = \frac{1}{2} e^x (x \sin x - x \cos x + \cos x)$$

$$\int x e^x \cos x = \frac{1}{2} e^x (x \sin x + x \cos x - \sin x)$$

# 10.4 Sum & Series

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$$

$$\begin{split} &\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \, (-1 < x \le 1) \\ &\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \, (-1 \le x \le 1) \\ &\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \, (-\infty < x < \infty) \\ &\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \, (-\infty < x < \infty) \end{split}$$

# 10.5 Misc

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f$$

$$x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .
- BEST theorem: the number of eulerian circuits in a directed graph is  $|\det(L_{ww})| \cdot \prod_{v \in V} (\deg(v) - 1)!$ .
- Tutte's Matrix

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

- · Cayley's Formula
  - Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each  $\emph{labeled}$  vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of *labeled* forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .
- Erdős-Gallai theorem

A sequence of nonnegative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only

if 
$$d_1+\cdots+d_n$$
 is even and  $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$  holds

for every  $1 \le k \le n$ .

Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \geq \cdots \geq a_n$  and  $b_1, \ldots, b_n$ 

is bigraphic if and only if 
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$ 

holds for every  $1 \le k \le n$ .

Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,b_1),\ldots,(a_n,b_n)$  of nonnegative integer pairs with  $a_1\geq$ 

$$\cdots \geq a_n$$
 is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1)$ 

$$1) + \sum_{i=k+1}^n \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have

 $A=\#\{\mbox{lattice points in the interior}\}+\frac{\#\{\mbox{lattice points on the boundary}\}}{2}-1.$ 

Möbius inversion formula

- $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
- $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- - A portion of a sphere cut off by a plane.
  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$  $\cos \theta)^2/3$ .
  - Area =  $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- · Lagrange multiplier
  - Optimize  $f(x_1,\ldots,x_n)$  when k constraints  $g_i(x_1,\ldots,x_n)=0$ .
  - Lagrangian function
  - Lagrangian random  $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n).$  The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines

- Line 1 : 
$${m v}_1 = {m p}_1 + t_1 {m d}_1$$

- Line 2 : 
$$v_2 = p_2 + t_2 d_2$$

- $n = d_1 \times d_2$
- $n_1 = d_1 \times n$

- 
$$c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1$$

- 
$$c_2 = p_2 + rac{(oldsymbol{p}_1 - oldsymbol{p}_2) \cdot oldsymbol{n}_1}{oldsymbol{d}_2 \cdot oldsymbol{n}_1} oldsymbol{d}_2$$

• Bernoulli numbers 
$$B_0-1, B_1^{\pm}=\pm\frac{1}{2}, B_2=\frac{1}{6}, B_3=0$$

$$\sum_{j=0}^m {m+1 \choose j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{array}{l} S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n = \sum_{i=0}^n S(n,i)(x)_i \\ \bullet \text{ Pentagonal number theorem} \end{array}$$

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers 
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{k=0}^{k} (-1)^{j} {n+1 \choose k} (k+1-j)^{n}$$

# 10.6 Number

· Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039, 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

• Number of paritions of n:

Maximum number of divisors:

n | 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 •  $\frac{\binom{n}{n}}{\binom{2n}{n}}$  2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

• Number of ways to partition a set of 
$$n$$
 labeled elements:  $n \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid 10 \mid 11 \mid 12 \mid 13$   
•  $B_n \mid 2 \mid 5 \mid 15 \mid 52 \mid 203 \mid 877 \mid 4140 \mid 21147 \mid 115975 \mid 7e5 \mid 4e6 \mid 3e7 \mid 115975 \mid 7e5 \mid 4e6 \mid 4e7 \mid 115975 \mid 4e7 \mid 4$ 

• Fibonacci numbers:  $\frac{n}{F_n}$  1 1 2 3 4 5 31 45 88 1346269 1e9 1e18

