```
Contents
                                         5.11 Minimum Steiner Tree . . . 14
                                         5.12 Count Cycles . . . . . . . . . 15
  Basic
                                      6 Math
  1.1 Default Code . . . . . . .
                                         6.1 Extended Euclidean Algo-
  1.2 .vimrc . . . . . . . . . . .
      Fast IO . . . . . . . . . . .
                                              rithm . . . . . . . . . . . . . . 15
  13
  1.4 Random . . . . . . . . .
                                            Floor & Ceil . . . . . . . . . . 15
  1.5 PBDS Tree . . . . . . . . .
                                         6.3 Legendre . . . . . . . . . . . . . . . 15
      Pragma . . . . . . . . . . .
                                         6.4 Simplex . . . . . . . . . . . . . . . . 15
                                         6.5 Simplex Construction . . . 16
2 Data Structure
                                         6.6 DiscreteLog . . . . . . . . 16
  2.1 Heavy-Light Decomposition
                                         6.7 Miller Rabin & Pollard Rho . 16
  2.2 Link Cut Tree . . . . . . . .
                                         6.8 XOR Basis . . . . . . . . . 16
  2.3 Treap . . . . . . . . . . . . .
                                         6.9 Linear Equation . . . . . . 17
  2.4 KD Tree . . . . . . . . . . . . .
                                         6.10 Chinese Remainder Theorem 17
  2.5 Leftist Tree . . . . . . . .
  2.6 Convex 1D/1D . . . . . . .
                                         6.11 Sart Decomposition . . . . 17
                                         Flow & Matching
   3.1 Dinic
  7 Polynomial
                                         7.1 FWHT . . . . . . . . . . . . . . . . 17
      MCMF . . . . . . . . . . . . .
  3.3
  3.4 Min Cost Circulation . . . .
                                         7.2 FFT . . . . . . . . . . . . . . . 18
  3.5 Gomory Hu . . . . . . . . .
                                         7.3 NTT . . . . . . . . . . . . . 18
  3.6 Stoer Wagner Algorithm . .
                                         7.4 Polynomial Operation . . . 18
  3.7 Bipartite Matching . . . . .
                                         7.5 Generating Function . . . . 20
  3.8 Kuhn Munkres Algorithm
                                  6
                                              Ordinary Generating Func-
  3.9 Max Simple Graph Matching
                                  6
  tion . . . . . . . . 20
                                              Exponential Generating
                                                   Function . . . . . . 20
  Geometry
                                         7.6 Bostan Mori . . . . . . . . 20
  4.1 Geometry Template . . . .
       Polar Angle Comparator . .
                                      8 String
      Minkowski Sum . . . . .
                                  8
                                         8.1 KMP Algorithm . . . . . . . 21
       Intersection of Circle and
  8.2 Manacher Algorithm . . . . 21
                                         8.3 Lyndon Factorization . . . . 21
      Tangent Line of Circles . . .
  4.6
                                         8.4 Suffix Array . . . . . . . . . 21
  4.7 Intersection of Line and
                                         8.5 Suffix Automaton . . . . . 21
       Convex Polygon . . . . . .
                                         8.6 Z-value Algorithm . . . . . 22
  4.8 Intersection of Line and
                                         8.7 Main Lorentz . . . . . . . . . 22
       Circle . . . . . . . . . . . .
  4.9 Point in Circle . . . . . .
                                         8.8 AC Automaton . . . . . . . 22
  4.10 Point in Convex . . . . . .
                                         8.9 Palindrome Automaton . . 22
  4.11 Half Plane Intersection
  4.12 Minimum Enclosing Circle .
                                      9 Misc
  9.1 Cyclic Ternary Search . . . 23
                                         9.2 Matroid . . . . . . . . . . . 23
  4.15 Delaunay Triangulation . . 10
  4.16 Voronoi Diagram . . . . . . 10
                                         9.3 Simulate Annealing . . . . 23
  4.17 Polygon Union . . . . . .
                                         9.4 Binary Search On Fraction . 23
  4.18 Tangent Point to Convex Hull 11
                                         9.5 Min Plus Convolution . . . 23
  4.19 Heart . . . . . . . . . . . . . . . . 11
  4.20 Rotating Sweep Line . . . . 11
  10 Notes
                                         10.1 Geometry . . . . . . . . . 23
  4.23 Calculate Points in Triangle 12
                                              Rotation Matrix . . . . . . 23
                                              Triangles . . . . . . . . . . 23
  Graph
                                              Quadrilaterals . . . . . . 24
  Spherical coordinates . . . 24

      5.2
      SCC
      12

      5.3
      2-SAT
      12

      5.4
      Dominator Tree
      12

                                              Green's Theorem . . . . . . 24
                                              Point-Line Duality . . . . . 24
  10.2 Trigonometry . . . . . . . 24
  5.6 Fast DMST . . . . . . . . . . . 13
                                         10.3 Calculus . . . . . . . . . 24
  10.4 Sum & Series . . . . . . . . 24
  5.9 Number of Maximal Clique 14
5.10 Minimum Mean Cycle . . . 14
                                         10.5 Misc . . . . . . . . . . . 24
                                         10.6 Number . . . . . . . . . . 25
      Basic
1.1 Default Code [2d2229]
#include <bits/stdc++.h>
using namespace std;
```

```
#include <bits/stdc++.h>
using namespace std;

#define iter(v) v.begin(),v.end()
#define SZ(v) int(v.size())
#define pb emplace_back
#define ff first
#define ss second

using ll = long long;
using pii = pair<int, int>;
using pll = pair<ll, ll>;

#ifdef zisk
void debug(){cerr << "\n";}
template<class T, class ... U>
```

```
void debug(T a, U ... b){cerr << a << " ", debug(b...)</pre>
template < class T> void pary(T 1, T r){
  while (1 != r) cerr << *1 << "
  cerr << "\n";</pre>
#else
#define debug(...) void()
#define pary(...) void()
#endif
template < class A, class B>
ostream& operator<<(ostream& o, pair<A,B> p) {    return o << '(' << p.ff << ',' << p.ss <<
                                  << p.ss << ')'; }
int main(){
  ios_base::sync_with_stdio(false); cin.tie(0);
1.2 .vimrc [476d29]
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
map <F9> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -
    Wextra -Wshadow -O2 -Dzisk -g -fsanitize=undefined,
    address<CR>
map <F8> :!./%:r<CR>
inoremap {<CR> {<CR>}<ESC>ko
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
inoremap fj <ESC>
vnoremap fj <ESC>
" -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG
1.3 Fast IO [4f6f0e]
char readchar() {
  const int N = 1 << 20;
  static char buf[N];
  static char *p = buf , *end = buf;
  if(p == end) {
    if((end = buf + fread(buf , 1 , N , stdin)) == buf)
         return EOF;
    p = buf;
  return *p++;
}
const int buf size = 524288;
struct Writer {
  char buf[buf_size]; int size = 0, ret;
  void flush() { ret = write(1, buf, size); size = 0; }
  void _flush(int sz) { if (sz + size > buf_size) flush
      (); }
  void write_char(char c) { _flush(1); buf[size++] = c;
  void write_int(int x) {
    const int len = 20;
     _flush(len);    <mark>int</mark> ptr = 0;
    if (x < 0) buf[size++] = '-', x = -x;</pre>
    if (x == 0) buf[size + (ptr++)] = '0';
    else for (; x; x /= 10) buf[size + (ptr++)] = '0' +
         x % 10;
    reverse(buf + size, buf + size + ptr);
    size += ptr:
}; // remember to call flush
1.4 Random [4cf9ed]
mt19937 rng(chrono::system clock::now().
    time_since_epoch().count());
1.5 PBDS Tree [7c702a]
#include <bits/extc++.h>
using namespace __gnu_pbds;
using Tree = tree<int, null_type, less<>, rb_tree_tag,
    tree_order_node_statistics_update>;
// .find_by_order(x)
// .order_of_key(x)
```

1.6 Pragma [6006f6]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,arch=skylake")
 __builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
1.7 SVG Writer [7adcc8]
class SVG {
  void p(string_view s) { o << s; }</pre>
  void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
     o << s.substr(0, i) << v, p(s.substr(i + 1), vs...)</pre>
  ofstream o; string c = "red";
public: // SVG svg("test.svg", 0, 0, 100, 100)
  SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f
     p("<svg xmlns='http://www.w3.org/2000/svg' "
        "viewBox='$ $ $ '>\n"
        "<style>*{stroke-width:0.5%;}</style>\n",
  x1, -y2, x2 - x1, y2 - y1); } ~SVG() { p("</svg>\n"); }
  void color(string nc) { c = nc; }
  void line(auto x1, auto y1, auto x2, auto y2) {
  p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n
        x1, -y1, x2, -y2, c); }
  void circle(auto x, auto y, auto r) {
 p("<circle cx='$' cy='$' r='$' stroke='$' "
    "fill='none'/>\n", x, -y, r, c); }
void text(auto x, auto y, string s, int w = 12) {
    "("<tay* x='$' y='$' fort circ='$ny'\$' (tox*t) b'</pre>
```

2 Data Structure

};

x, -y, w, s); }

2.1 Heavy-Light Decomposition [f2dbca]

p("<text x='\$' y='\$' font-size='\$px'>\$</text>\n",

```
struct HLD{ // 1-based
  int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
  explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1),
  ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
   }
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  void build(){ dfs(1, 1), cut(1, 1); }
int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // Lca
};
```

2.2 Link Cut Tree [cf4f34]

```
// 1-based
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node(): pa{0}, ch{0, 0}, rev{false}, size{1}, v{},
      sum{}, rsum{}, sv{}, sub{}, vir{} {}
  };
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is_rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false:
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
    cur.rev ^= 1;
  void rotate(int u) {
    int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[l ^ 1]) o[cur.ch[l ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  void splay(int u) {
    vector<int> stk = {u};
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is_root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    up(u);
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
      cur.vir = cur.vir + o[rc].sub - o[last].sub;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0;
    for (access(u); u; u = lc) down(la = u);
    return la;
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); }
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
    return find_root(u) == find_root(v); }
```

```
void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
#undef rc
};
2.3 Treap [2ac37e]
mt19937 rng(880301);
struct node {
  11 data; int sz;
  node *1, *r:
  node(ll k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
}:
node pool[1000010]; int pool_cnt = 0;
node *newnode(ll k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
           a;
  return b->down(), b->l = merge(a, b->l), b->up(), b;
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o \rightarrow r, a \rightarrow r, b, k), <math>a \rightarrow up();
  else b = o, split(o->1, a, b->1, k), b->up();
// a: size k, b: size n - k
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();</pre>
  o->down();
  if (sz(o->1) + 1 <= k)
   a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
node *kth(node *o, ll k) { // 1-based
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, 11 key) { // num of key < key</pre>
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, 11 k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, ll k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
tuple<node*, node*, node*> interval(node *&o, int 1,
    int r) { // 1-based
  node *a, *b, *c; // b: [l, r]
```

```
split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
   return make_tuple(a, b, c);
}
2.4 KD Tree [375ca2]
 namespace kdt {
   int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
   yl[maxn], yr[maxn];
   point p[maxn];
   int build(int 1, int r, int dep = 0) {
     if (1 == r) return -1;
     function<bool(const point &, const point &)> f =
       [dep](const point &a, const point &b) {
  if (dep & 1) return a.x < b.x;</pre>
         else return a.y < b.y;</pre>
       };
     int m = (1 + r) >> 1;
     nth_element(p + 1, p + m, p + r, f);
     x1[m] = xr[m] = p[m].x;
     yl[m] = yr[m] = p[m].y;
     lc[m] = build(1, m, dep + 1);
     if (~lc[m]) {
       xl[m] = min(xl[m], xl[lc[m]]);
       xr[m] = max(xr[m], xr[lc[m]]);
       yl[m] = min(yl[m], yl[lc[m]]);
       yr[m] = max(yr[m], yr[lc[m]]);
     rc[m] = build(m + 1, r, dep + 1);
     if (~rc[m]) {
       x1[m] = min(x1[m], x1[rc[m]]);
       xr[m] = max(xr[m], xr[rc[m]]);
       yl[m] = min(yl[m], yl[rc[m]]);
       yr[m] = max(yr[m], yr[rc[m]]);
     return m;
   bool bound(const point &q, int o, long long d) {
     double ds = sqrt(d + 1.0);
     if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
         q.y < yl[o] - ds || q.y > yr[o] + ds
       return false;
     return true;
   long long dist(const point &a, const point &b) {
     return (a.x - b.x) * 111 * (a.x - b.x) + (a.y - b.y) * 111 * (a.y - b.y);
   void dfs(
       const point &q, long long &d, int o, int dep = 0)
     if (!bound(q, o, d)) return;
     long long cd = dist(p[o], q);
     if (cd != 0) d = min(d, cd);
     if ((dep & 1) && q.x < p[o].x ||</pre>
         !(dep & 1) && q.y < p[o].y) {
       if (~lc[o]) dfs(q, d, lc[o], dep + 1);
       if (~rc[o]) dfs(q, d, rc[o], dep + 1);
     } else {
       if (~rc[o]) dfs(q, d, rc[o], dep + 1);
       if (~lc[o]) dfs(q, d, lc[o], dep + 1);
     }
   }
   void init(const vector<point> &v) {
     for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
     root = build(0, v.size());
   long long nearest(const point &q) {
     long long res = 1e18;
     dfs(q, res, root);
     return res;
} // namespace kdt
 2.5 Leftist Tree [e91538]
 struct node {
  11 v, data, sz, sum;
node *1, *r;
   node(11 k)
     : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
ll sz(node *p) { return p ? p->sz : 0; }
```

```
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->1)) swap(a->r, a->1);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
 return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
```

2.6 Convex 1D/1D [a449dd]

```
template < class T>
struct DynamicHull {
  struct seg { int x, l, r; };
  T f; int C; deque<seg> dq; // range: 1~C
explicit DynamicHull(T _f, int _C): f(_f), C(_C) {}
  // max t s.t. f(x, t) \Rightarrow f(y, t), x < y, maintain max
  int intersect(int x, int y) {
    int l = 0, r = C + 1;
    while (1 + 1 < r) {
      int mid = (1 + r) / 2;
      if (f(x, mid) >= f(y, mid)) l = mid;
      else r = mid;
    }
    return 1;
  void push_back(int x) {
    for (int i; !dq.empty() &&
        (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back(). r + 1 <= C) dq.pb(seg({x, dq.back().</pre>
        r + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
 }
};
```

Flow & Matching

3.1 Dinic [801a71]

```
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
   ll to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
   n = n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
   g[u].pb(edge{v, cap, 0, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
   dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
```

```
int cur = q.front(); q.pop();
      for (auto &e : g[cur])
        if (dis[e.to] == -1 && e.flow != e.cap) {
          q.push(e.to);
          dis[e.to] = dis[cur] + 1;
      }
    }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
          g[e.to][e.rev].flow -= df;
          return df;
        }
      }
    dis[u] = -1;
    return 0;
  11 maxflow(int _s, int _t) {
    s = _s; t = _t;
    11 \text{ flow} = 0, \text{ df};
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
};
      Bounded Flow [758826]
3.2
struct BoundedFlow : Dinic {
  vector<ll> tot;
  void init(int _n) {
    Dinic::init(_n + 2);
    tot.assign(n, 0);
  void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
    g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
  bool feasible() {
    11 \text{ sum} = 0;
    int vs = n - 2, vt = n - 1;
    for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
        add_edge(vs, i, 0, tot[i]), sum += tot[i];
      else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
        g[vs].pop_back(), g[i].pop_back();
      else if(tot[i] < 0)</pre>
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
    add_edge(_t, _s, 0, INF);
    if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    //return x + maxflow(_s, _t); // max
 }
};
3.3 MCMF [8d3644]
struct MCMF { // 0-based, O(SPFA * |f|)
  struct edge {
    11 from, to, cap, flow, cost, rev;
  int n;
```

int s, t; 11 mx;

```
//mx: maximum amount of flow
  vector<vector<edge>> g;
  vector<ll> dis, up;
  bool BellmanFord(ll &flow, ll &cost) {
    vector<edge*> past(n);
    vector<int> inq(n);
    dis.assign(n, INF); up.assign(n, 0);
    queue<int> q;
    q.push(s), inq[s] = 1;
    up[s] = mx - flow, past[s] = 0, dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      if (!up[u]) continue;
      for (auto &e : g[u])
        if (e.flow != e.cap &&
            dis[e.to] > dis[u] + e.cost) {
          dis[e.to] = dis[u] + e.cost, past[e.to] = &e;
          up[e.to] = min(up[u], e.cap - e.flow);
          if (!inq[e.to]) inq[e.to] = 1, q.push(e.to);
    if (dis[t] == INF) return 0;
    flow += up[t], cost += up[t] * dis[t];
    for (ll i = t; past[i]; i = past[i]->from) {
      auto &e = *past[i];
      e.flow += up[t], g[e.to][e.rev].flow -= up[t];
   }
    return 1:
  pll MinCostMaxFlow(int _s, int _t) {
    s = _s, t = _t;
11 flow = 0, cost = 0;
    while (BellmanFord(flow, cost));
    return pll(flow, cost);
  void init(int _n, ll _mx) {
   n = _n, mx = _mx;
    g.assign(n, vector<edge>());
  void add_edge(int a, int b, ll cap, ll cost) {
    g[a].pb(edge{a, b, cap, 0, cost, SZ(g[b])});
    g[b].pb(edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
};
3.4 Min Cost Circulation [47cf18]
struct MinCostCirculation { // 0-based, O(VE * ELogC)
  struct edge {
   11 from, to, cap, fcap, flow, cost, rev;
  int n;
  vector<edge*> past;
  vector<vector<edge>> g;
  vector<ll> dis;
  void BellmanFord(int s) {
    vector<int> inq(n);
    dis.assign(n, INF);
    queue<int> q;
    auto relax = [&](int u, ll d, edge *e) {
      if (dis[u] > d) {
        dis[u] = d, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
     }
    };
    relax(s, 0, 0);
```

while (!q.empty()) {

}

}

int u = q.front();

q.pop(), inq[u] = 0;

for (auto &e : g[u])
if (e.cap > e.flow)

void try_edge(edge &cur) {

auto &e = *past[i];

BellmanFord(cur.to);

) {

relax(e.to, dis[u] + e.cost, &e);

if (dis[cur.from] + cur.cost < 0) {</pre>

++cur.flow, --g[cur.to][cur.rev].flow;

if (cur.cap > cur.flow) return ++cur.cap, void();

for (int i = cur.from; past[i]; i = past[i]->from

};

```
++e.flow, --g[e.to][e.rev].flow;
      }
    ++cur.cap;
  void solve(int mxlg) { // mxlg >= log(max cap)
    for (int b = mxlg; b >= 0; --b) {
  for (int i = 0; i < n; ++i)</pre>
         for (auto &e : g[i])
      e.cap *= 2, e.flow *= 2;
for (int i = 0; i < n; ++i)
        for (auto &e : g[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  }
  void init(int _n) {
    n = n;
    past.assign(n, nullptr);
    g.assign(n, vector<edge>());
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a)}
         == b));
    g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1});
  }
};
3.5
      Gomory Hu [82d968]
void GomoryHu(Dinic &flow) { // 0-based
  int n = flow.n;
  vector<int> par(n);
  for (int i = 1; i < n; ++i) {</pre>
    flow.reset();
    add_edge(i, par[i], flow.maxflow(i, par[i]));
    for (int j = i + 1; j < n; ++j)
  if (par[j] == par[i] && ~flow.dis[j])</pre>
        par[j] = i;
}
3.6 Stoer Wagner Algorithm [a9917b]
struct StoerWagner { // 0-based, O(V^3)
  int n;
  vector<int> vis, del;
  vector<ll> wei;
  vector<vector<ll>> edge:
  void init(int _n) {
    n = _n;
    del.assign(n, 0);
    edge.assign(n, vector<11>(n));
  void add_edge(int u, int v, ll w) {
    edge[u][v] += w, edge[v][u] += w;
  void search(int &s, int &t) {
    vis.assign(n, 0); wei.assign(n, 0);
    s = t = -1;
    while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
         if (!del[i] && !vis[i] && mx < wei[i])</pre>
           cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)
        if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
    }
  11 solve() {
    11 ret = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
      for (int j = 0; j < n; ++j)</pre>
         edge[x][j] = (edge[j][x] += edge[y][j]);
    return ret;
```

3.7 Bipartite Matching [013c49]

```
//min vertex cover: take all unmatched vertices in L
    and find alternating tree,
//ans is not reached in L + reached in R
// O(VE)
int n; // 1-based, max matching
int mx[maxn], my[maxn];
bool adj[maxn][maxn], vis[maxn];
bool dfs(int u) {
  if (vis[u]) return 0;
  vis[u] = 1;
  for (int v = 1; v <= n; v++) {
    if (!adj[u][v]) continue;
    if (!my[v] || (my[v] && dfs(my[v]))) {
     mx[u] = v, my[v] = u;
      return 1;
   }
  }
  return 0;
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!\sim my[e] || (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
for (int i = 0; i < nl; ++i)</pre>
      if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
        }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    }
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
  nl = _nl, nr = _nr;
    g.assign(nl, vector<int>());
};
     Kuhn Munkres Algorithm [683e0a]
```

```
struct KM { // 0-based, maximum matching, O(V^3)
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll>> hl, hr, slk;
  vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    n = _n;
    // -INF for perfect matching
    w.assign(n, vector<ll>(n, 0));
    pre.assign(n, 0);
```

```
qu.assign(n, 0);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
        0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!v1[x] \&\& s1k[x] >= (d = h1[x] + hr[y] -
               w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !slk[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
         ; hr.assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
     hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  }
};
```

3.9 Max Simple Graph Matching [907d7c]

```
struct Matching { // 0-based, 0(V^3)
 queue<int> q; int n;
 vector<int> fa, s, vis, pre, match;
 vector<vector<int>> g;
 int Find(int u)
 { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
 int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
     if (vis[x] == tk) return x;
      vis[x] = tk;
     x = Find(pre[match[x]]);
 void Blossom(int x, int y, int 1) {
   for (; Find(x) != 1; x = pre[y]) {
     pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
 bool Bfs(int r) {
   iota(iter(fa), 0); fill(iter(s), -1);
    q = queue < int > (); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : g[x])
        if (s[u] == -1) {
          if (pre[u] = x, s[u] = 1, match[u] == n) {
            for (int a = u, b = x, last;
    b != n; a = last, b = pre[a])
              last = match[b], match[b] = a, match[a] =
                   b;
            return true;
```

```
q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int 1 = LCA(u, x);
          Blossom(x, u, 1); Blossom(u, x, 1);
    return false;
  Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis(n)
  + 1), pre(n + 1, n), match(n + 1, n), g(n) {} void add_edge(int u, int v)
  { g[u].pb(v), g[v].pb(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

3.10 Stable Marriage

```
1: Initialize m \in M and w \in W to free
 2: while \exists free man m who has a woman w to propose to do
        w \leftarrow \text{first woman on } m \text{'s list to whom } m \text{ has not yet proposed}
3:
 4:
        if \exists some pair (m', w) then
 5:
             if w prefers m to m' then
 6:
                 m' \leftarrow free
7:
                 (m,w) \leftarrow \textit{engaged}
 8:
             end if
9:
         else
10:
             (m,w) \leftarrow \mathsf{engaged}
11:
         end if
```

12: end while 3.11 Flow Model

- · Maximum/Minimum flow with lower bound / Circulation problem
- 1. Construct super source S and sink T.
- 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v)>0, connect S o v with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \, o \, s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f'
 eq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
- 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x \rightarrow y$ with (cost,cap) = (c,1) if c>0, otherwise connect $y\to x$ with (cost, cap)=(-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=(0, d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer T
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v$, $v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity
- 2w(v)
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. Let $w'(u,v) = w(u,v) \mu(u) \mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 2. Find the minimum weight matching M with w' . The answer is $\sum \mu(v) +$ w'(M).

- · Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
 - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers

$$\min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$
 $p_{u} \geq 0$

Geometry

/ abs2(dp);

Geometry Template [86f0f1]

```
using ld = 11;
                                                             using pdd = pair<ld, ld>;
                                                             #define X first
                                                             #define Y second
                                                             // Ld eps = 1e-7;
                                                             pdd operator+(pdd a, pdd b)
                                                             { return {a.X + b.X, a.Y + b.Y}; }
                                                             pdd operator-(pdd a, pdd b)
                                                             { return {a.X - b.X, a.Y - b.Y}; }
                                                             pdd operator*(ld i, pdd v)
                                                             { return {i * v.X, i * v.Y}; }
                                                             pdd operator*(pdd v, ld i)
                                                             { return {i * v.X, i * v.Y}; }
                                                             pdd operator/(pdd v, ld i)
                                                             { return {v.X / i, v.Y / i}; }
                                                             ld dot(pdd a, pdd b)
                                                             { return a.X * b.X + a.Y * b.Y; }
                                                             ld cross(pdd a, pdd b)
                                                             { return a.X * b.Y - a.Y * b.X; }
                                                             ld abs2(pdd v)
                                                             { return v.X * v.X + v.Y * v.Y; };
                                                             ld abs(pdd v)
                                                             { return sqrt(abs2(v)); };
                                                             int sgn(ld v)
                                                             { return v > 0 ? 1 : (v < 0 ? -1 : 0); }
                                                             // int sgn(ld v){    return v > eps ? 1 : ( v < -eps ? -1
                                                                  : 0);
                                                             int ori(pdd a, pdd b, pdd c)
                                                             { return sgn(cross(b - a, c - a)); }
                                                             bool collinearity(pdd a, pdd b, pdd c)
                                                             { return ori(a, b, c) == 0; }
                                                             bool btw(pdd p, pdd a, pdd b)
                                                             { return collinearity(p, a, b) && sgn(dot(a - p, b - p)
                                                                  ) <= 0: }
                                                             bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
                                                               if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
                                                                    p2) || btw(p4, p1, p2))
                                                                  return true;
                                                               return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&
  ori(p3, p4, p1) * ori(p3, p4, p2) < 0;</pre>
                                                             pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
                                                               ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
                                                               return (p4 * a123 - p3 * a124) / (a123 - a124);
                                                             pdd perp(pdd p1)
                                                             { return pdd(-p1.Y, p1.X); }
                                                             pdd projection(pdd p1, pdd p2, pdd p3)
                                                             { return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
                                                                 p2 - p1); }
                                                             pdd reflection(pdd p1, pdd p2, pdd p3)
5. For v \in G, connect it with sink v \to t with capacity K+2T-(\sum_{e \in E(v)} u \not| (q)) return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) / (q)
                                                                   abs2(p2 - p1) * 2; }
                                                             pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
                                                                  , pdd r) {
                                                               pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
                                                                    dot(dp, dq));
                                                               return q0 + pdd(cross(r - p0, num), dot(r - p0, num))
```

```
} // from line p0--p1 to q0--q1, apply to r
```

4.2 Polar Angle Comparator [808e89]

4.3 Minkowski Sum [98abff]

```
void reorder_poly(vector<pdd>& pnts){
  int mn = 0;
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn</pre>
        ].Y && pnts[i].X < pnts[mn].X))</pre>
  rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
  reorder_poly(P);
  reorder_poly(Q);
  int psz = P.size();
  int qsz = Q.size();
  P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
  vector<pdd> ans;
 int i = 0, j = 0;
while(i < psz || j < qsz){</pre>
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1] - P[i], Q[j + 1] - Q[j])
    if(t >= 0) i++;
    if(t <= 0) j++;</pre>
  return ans;
}
```

4.4 Intersection of Circle and Convex Polygon

```
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
 if(a > r){
   S = (C/2)*r*r;
    h = a*b*sin(C)/c;
   if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
   theta = PI - B - asin(sin(B)/r*a);
   S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double areaPolyCircle(const vector<pdd> poly,const pdd
    &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=\_area(poly[i]-0,poly[(i+1)\%SZ(poly)]-0,r)*ori(0,
        poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
}
```

4.5 Intersection of Circles [f7a2fe]

```
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
  double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2), d =
        sqrt(d2);
```

4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
    int sign1 ){
  vector<Line> ret;
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
        v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sgn(p1.X - p2.X) == 0 and
        sgn(p1.Y - p2.Y) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret;
```

4.7 Intersection of Line and Convex Polygon

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 \mid | cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int 1, int r) {
    for (int t = 1; (1 + 1) % n != r; ) {
      int m = ((1 + r + (1 < r? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(t) ? 1 : r) = m;
    return (1 + !cmpL(r)) % n;
  };
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned in the same order as the line hits the
      convex */
  return res;
} // convex cut: (r, l]
```

4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
    a, pdd b) {
    pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
    ;
    double s = cross(b - a, c - a), h2 = r * r - s * s /
        abs2(b - a);
```

```
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  if (sgn(h2) < 0) return {};</pre>
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return \{p - h, p + h\};
4.9 Point in Circle [ecf954]
// return q's relation with circumcircle of tri(p[0],p
    [1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
   int128 det = 0;
  for (int i = 0; i < 3; ++i)
   }
  return det > 0; // in: >0, on: =0, out: <0
}
4.10 Point in Convex [f86640]
bool PointInConvex(const vector<pll> &C, pll p, bool
    strict = true) {
  int a = 1, b = SZ(C) - 1, r = !strict;
 if (SZ(C) == 0) return false;
if (SZ(C) < 3) return r && btw(C[0], C.back(), p);</pre>
  if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
  if (ori(C[0], C[a], p) >= r || ori(C[0], C[b], p) <=</pre>
      -r)
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (ori(C[0], C[c], p) > 0 ? b : a) = c;
  return ori(C[a], C[b], p) < r;</pre>
}
4.11 Half Plane Intersection [dfb833]
// from 8BQube
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
    .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0; // C^4
/* Having solution, check size > 2 */
```

```
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(iter(arr), [&](Line a, Line b) -> int {
    if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
      return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
         -1)
      continue;
    while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.
        back()))
      dq.pop_back();
    while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
   dq.pb(p);
 while (SZ(dq) >= 3 \&\& !isin(dq[0], dq[SZ(dq) - 2], dq
      .back()))
    dq.pop_back();
  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
   dq.pop_front();
  return vector<Line>(iter(dq));
}
```

4.12 Minimum Enclosing Circle [5af6d5]

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
 random_shuffle(iter(pts));
```

```
pdd c = pts[0];
  1d r = 0;
  for(int i = 1; i < SZ(pts); i++){</pre>
    if(abs(pts[i] - c) <= r) continue;
c = pts[i]; r = 0;</pre>
    for(int j = 0; j < i; j++){</pre>
      if(abs(pts[j] - c) <= r) continue;</pre>
      c = (pts[i] + pts[j]) / 2;
      r = abs(pts[i] - c);
      for(int k = 0; k < j; k++){</pre>
        if(abs(pts[k] - c) > r)
          tie(c, r) = circumcenter(pts[i], pts[j], pts[
              k]);
  return {c, r};
4.13 3D Point [badbbd]
// Copy from 8BQube
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x
       (_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
    .y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
```

4.14 ConvexHull3D [156311]

Point u = axis / abs(axis);

```
struct convex_hull_3D {
struct Face {
 int a, b, c;
```

Point e1 = b - a; Point e2 = c - a;

e1 = e1 / abs(e1);

e2 = e2 / abs(e2);

Point p = u - a;

* s;

}

e2 = e2 - e1 * dot(e2, e1);

return pdd(dot(p, e1), dot(p, e2));

double s = sin(angle), c = cos(angle);

Point rotate_around(Point p, double angle, Point axis)

return u * dot(u, p) * (1 - c) + p * c + cross(u, p)

```
Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
                                                             struct Edge {
                                                               int id; // oidx[id]
}; // return the faces with pt indexes
                                                               list<Edge>::iterator twin;
vector<Face> res;
vector<Point> P;
                                                               Edge(int _id = 0):id(_id) {}
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
                                                             struct Delaunay { // 0-base
  int n = SZ(P);
                                                               int n;
  if (n <= 2) return; // be careful about edge case</pre>
                                                               vector<int> oidx:
  // ensure first 4 points are not coplanar
                                                               vector<list<Edge>> head; // result udir. graph
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn
                                                               vector<pll> p;
      (abs2(P[0] - p)) != 0; }));
                                                               Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
  swap(P[2], *find_if(iter(P), [&](auto p) { return sgn
                                                                   head(n), p(n) {
  (abs2(cross3(p, P[0], P[1]))) != 0; }));
swap(P[3], *find_if(iter(P), [&](auto p) { return sgn
                                                                 iota(iter(oidx), 0);
                                                                 for (int i = 0; i < n; ++i) head[i].clear();</pre>
      (volume(P[0], P[1], P[2], p)) != 0; }));
                                                                 sort(iter(oidx), [&](int a, int b)
  vector<vector<int>> flag(n, vector<int>(n));
                                                                     res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
                                                                 for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
  for (int i = 3; i < n; ++i) {</pre>
                                                                 divide(0, n - 1);
    vector<Face> next;
    for (auto f : res) {
                                                               void addEdge(int u, int v) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
                                                                 head[u].push_front(Edge(v));
                                                                 head[v].push_front(Edge(u));
      if (d <= 0) next.pb(f);</pre>
                                                                 head[u].begin()->twin = head[v].begin();
      int ff = (d > 0) - (d < 0);
                                                                 head[v].begin()->twin = head[u].begin();
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
                                                               void divide(int 1, int r) {
                                                                 if (1 == r) return;
                                                                 if (1 + 1 == r) return addEdge(1, 1 + 1);
    for (auto f : res) {
                                                                 int mid = (1 + r) \gg 1, nw[2] = \{1, r\};
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
                                                                 divide(l, mid), divide(mid + 1, r);
                                                                 auto gao = [&](int t) {
          next.emplace_back(x, y, i);
                                                                   pll pt[2] = {p[nw[0]], p[nw[1]]};
                                                                   for (auto it : head[nw[t]]) {
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
                                                                     int v = ori(pt[1], pt[0], p[it.id]);
    res = next;
                                                                     if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
                                                                         id]) < abs2(pt[1] - pt[0])))
 }
                                                                       return nw[t] = it.id, true;
bool same(Face s, Face t) {
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
                                                                   return false;
       return 0:
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
                                                                 while (gao(0) || gao(1));
                                                                 addEdge(nw[0], nw[1]); // add tangent
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
                                                                 while (true) {
       return 0;
                                                                   pll pt[2] = {p[nw[0]], p[nw[1]]};
                                                                   int ch = -1, sd = 0;
for (int t = 0; t < 2; ++t)</pre>
  return 1:
                                                                     for (auto it : head[nw[t]])
int polygon_face_num() {
  int ans = 0;
                                                                       if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
  for (int i = 0; i < SZ(res); ++i)</pre>
                                                                             -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
    ans += none_of(res.begin(), res.begin() + i, [&](
                                                                            id])))
        Face g) { return same(res[i], g); });
                                                                         ch = it.id, sd = t;
  return ans;
                                                                   if (ch == -1) break; // upper common tangent
                                                                   for (auto it = head[nw[sd]].begin(); it != head[
double get_volume() {
                                                                       nw[sd]].end(); )
  double ans = 0;
                                                                     if (seg_strict_intersect(pt[sd], p[it->id], pt[
  for (auto f : res)
                                                                         sd ^ 1], p[ch]))
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
                                                                       head[it->id].erase(it->twin), head[nw[sd]].
        ]);
                                                                           erase(it++);
                                                                     else ++it;
  return fabs(ans / 6);
                                                                   nw[sd] = ch, addEdge(nw[0], nw[1]);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
                                                            }:
      .z) * (p3.y - p1.y);
                                                                   Voronoi Diagram [e4f408]
                                                             4.16
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
      .x) * (p3.z - p1.z);
                                                             // all coord. is even, you may want to call
  double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
                                                                 halfPlaneInter after then
      .y) * (p3.x - p1.x);
                                                             vector<vector<Line>> vec;
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
                                                             void build_voronoi_line(int n, vector<pll> &pts) {
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
                                                               Delaunay tool(n, pts); // Delaunay
       * a + b * b + c * c);
                                                               vec.clear(), vec.resize(n);
                                                               for (int i = 0; i < n; ++i)</pre>
                                                                 for (auto e : tool.head[i]) {
// n^2 delaunay: facets with negative z normal of // convexhull of (x, y, x^2 + y^2), use a pseudo-point
                                                                   int u = tool.oidx[i], v = tool.oidx[e.id];
                                                                   pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v])
// (0, 0, inf) to avoid degenerate case
                                                                        - pts[u]);
4.15 Delaunay Triangulation [6a9916]
                                                                   vec[u].pb(Line(m, m + d));
/* Delaunay Triangulation:
   Given a sets of points on 2D plane, find a
```

4.17 Polygon Union [9fbf66]

triangulation such that no points will strictly

inside circumcircle of any triangle. */

```
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                                                 fruit_advantages
// from 8BQube
                                                              pll d; int u, v;
ld rat(pll a, pll b) {
  return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
 // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
  1d res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
  pll C = q[b], D = q[(b + 1) % SZ(q)];</pre>
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            1d sa = cross(D - C, A - C), sb = cross(D - C)
                C, B - C);
            segs.pb(sa / (sa - sb), sgn(sc - sd));
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D
               - C)) > 0) {
            segs.pb(rat(C - A, B - A), 1);
            segs.pb(rat(D - A, B - A), -1);
        }
      }
      sort(iter(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
                                                            }
      res += cross(A, B) * sum;
  return res / 2;
                                                                 a-c
4.18 Tangent Point to Convex Hull [523bc1]
                                                                  strict:
// from 8BOube
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  };
                                                            }
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
4.19 Heart [082d19]
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
                                                                });
    (center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
                                                                ]; };
  double m = 2. * (x1 * y2 - y1 * x2);
  pdd center;
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
      y1 - y2)) / m;
  center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
      y2 * y2) / m;
  return center + p0;
pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
     / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
                                                                  });
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
4.20 Rotating Sweep Line [f5f689]
struct Event {
```

```
bool operator<(const Event &b) const {</pre>
    int ret = cmp(d, b.d, false);
    return ret == -1 ? false : ret; } // no tie-break
void rotatingSweepLine(const vector<pll> &p) {
  const int n = SZ(p);
  vector<Event> e; e.reserve(n * (n - 1));
  for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++) // pos[i] < pos[j] when
         the event occurs
      if (i != j) e.pb(p[j] - p[i], i, j);
  sort(iter(e));
  vector<int> ord(n), pos(n);
  iota(iter(ord), 0);
  sort(iter(ord), [&](int i, int j) { // initial order
      return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
          X < p[j].X; \});
  for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
  // initialize
  for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
    // do something
    vector<pii> tmp;
    for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
      tmp.pb(pii(e[j].u, e[j].v));
    sort(iter(tmp), [&](pii x, pii y){
        return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
             ], pos[y.ss]); });
    for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
  tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
        make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
            pos[x]);
4.21 Vector In Poly [c6d0fa]
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
  return ori(a, b, p) >= strict && ori(a, p, c) >=
// whether vector{cur, p} in counter-clockwise order
    prv, cur, nxt
bool inside(pll prv, pll cur, pll nxt, pll p, int
    strict) {
  if (ori(cur, nxt, prv) >= 0)
    return btwangle(cur, nxt, prv, p, strict);
  return !btwangle(cur, prv, nxt, p, !strict);
4.22 Convex Hull DP [92fd4b]
sort(iter(pts), [&](pll x, pll y) {
    return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;</pre>
auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
// DP for convex hull vertices (no points on edges)
auto solve = [\&](int bottom) { // <math>O(n^3)
  pll 0 = pts[bottom];
  vector<pii> trans;
  for (int j = bottom + 1; j < n; j++)
    for (int k = bottom + 1; k < n; k++) {</pre>
      if (ori(0, pts[j], pts[k]) <= 0) continue;</pre>
      // check whether j->k is legal
      trans.pb(pii(j, k));
  sort(iter(trans), [&](pii x, pii y) -> bool{
      int tmp = cmp(getvec(x), getvec(y), false);
      if (tmp != -1) return tmp;
      pll v = getvec(x);
      return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether bottom -> j is legal
    // init trans -> j
  for (auto [i, j] : trans) {
    // normal trans i -> j
```

```
National Taiwan University
                                                  fruit_advantages
  for (int j = bottom + 1; j < n; j++) {</pre>
                                                                   if((now != par && low[now] >= in[par]) || (now ==
                                                                        par && SZ(g[now]) == 0)){
    // check whether j -> bottom is legal
    // end trans j ->
                                                                     make_bcc();
 }
                                                                     for(int v = stk.back();; v = stk.back()){
                                                                       stk.pop_back(), add_v(v);
for(int i = 0; i < n; i++) solve(i);</pre>
                                                                       for(int e : up[v]) add_e(e);
                                                                       if(v == now) break;
4.23 Calculate Points in Triangle [bf746f]
                                                                     if(now != par) add_v(par);
// all points are distinct
                                                                   }
// cnt[i][j] = # of point k s.t. strictly above ij, and
                                                                 };
                                                                 for(int i = 0; i < n; i++)</pre>
// cnt2[i][j] = # of points k s.t. strictly in ij
                                                                   if(in[i] == -1) _dfs(_dfs, i, i, -1);
// preprocess space: O(n^2), time: O(n^3), query time:
                                                               }
    0(1)
                                                            };
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
                                                             5.2
                                                                  SCC [2c9a01]
  for (int j = 0; j < n; j++){
    if (pts[i] >= pts[j]) continue;
                                                             struct SCC{ // 0-based, output reversed topo order
    for (int k = 0; k < n; k++) {
                                                               int n, cnt = 0;
      if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
                                                               vector<vector<int>> g;
        int tmp = ori(pts[i], pts[j], pts[k]);
                                                               vector<int> sccid;
        if (tmp > 0) cnt[i][j]++; // only for i < j
                                                               explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
void add_edge(int u, int v){
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
      }
                                                                 g[u].pb(v);
   }
                                                               void build(){
auto calc_tri = [&](array<int, 3> arr) { // strictly
                                                                 vector<int> in(n, -1), low(n), stk;
    inside
                                                                 vector<bool> instk(n);
  sort(iter(arr), [&](int x, int y){ return pts[x] <</pre>
                                                                 int ts = 0;
      pts[y]; });
                                                                 auto dfs1 = [&](auto dfs, int now) -> void{
  auto [x, y, z] = arr;
                                                                   stk.pb(now); instk[now] = true;
  int tmp = ori(pts[x], pts[y], pts[z]);
                                                                   in[now] = low[now] = ts++;
  if (tmp == 0) return 0;
                                                                   for(int i : g[now]){
  else if (tmp < 0)</pre>
                                                                     if(in[i] == -1)
    dfs(dfs, i), low[now] = min(low[now], low[i])
        y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
                                                                     else if(instk[i] && in[i] < in[now])</pre>
      x][z];
                                                                       low[now] = min(low[now], in[i]);
};
                                                                   if(low[now] == in[now]){
5
     Graph
                                                                     for(; stk.back() != now; stk.pop_back())
                                                                       sccid[stk.back()] = cnt, instk[stk.back()] =
5.1 BCC [d04ebe]
                                                                            false;
                                                                     sccid[now] = cnt++, instk[now] = false, stk.
struct BCC{ // O-based, allow multi edges but not allow
                                                                          pop_back();
     Loops
  int n, m, cnt = 0;
                                                                 };
  // n:|V|, m:|E|, cnt:#bcc
                                                                 for(int i = 0; i < n; i++)</pre>
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
                                                                   if(in[i] == -1) dfs1(dfs1, i);
  vector<vector<int>> bcc_v, bcc_e;
                                                               }
  vector<vector<pii>>> g; // original graph
                                                            };
  vector<pii> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
                                                            5.3
                                                                  2-SAT [0686a5]
    n(_n), m(SZ(_edges)), g(_n), edges(_edges){
      for(int i = 0; i < m; i++){
  auto [u, v] = edges[i];</pre>
                                                             struct SAT { // 0-based
                                                               int n:
        g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
                                                               vector<bool> istrue;
      }
                                                               SCC scc;
                                                               SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
  void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
                                                               int neg(int a) {
  // modify these if you need more information
void add_v(int v){ bcc_v.back().pb(v); }
                                                                 return a >= n ? a - n : a + n;
  void add_e(int e){ bcc_e.back().pb(e); }
                                                               void add_clause(int a, int b) {
  void build(){
                                                                 scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
    vector<int> in(n, -1), low(n, -1), stk;
    vector<vector<int>> up(n);
                                                               bool solve()
    int ts = 0;
                                                                 scc.build();
    auto _dfs = [&](auto dfs, int now, int par, int pe)
                                                                 for (int i = 0; i < n; ++i) {</pre>
         -> void{
```

if(pe != -1) up[now].pb(pe);

in[now] = low[now] = ts++;

for(auto [v, e] : g[now]){ if(e == pe) continue;

dfs(dfs, v, now, e);

if(in[v] < in[now]) up[now].pb(e);</pre>

low[now] = min(low[now], in[v]);

low[now] = min(low[now], low[v]);

if(in[v] != -1){

continue;

stk.pb(now);

5.4 Dominator Tree [2da9bb]

false;

return true;

}

};

```
struct Dominator {
 int n:
 vector<vector<int>> g, r, rdom; int tk;
```

istrue[i + n] = !istrue[i];

if (scc.sccid[i] == scc.sccid[i + n]) return

istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>

```
vector<int> dfn, rev, fa, sdom, dom, val, rp;
  Dominator(int _n) : n(_n), g(n), r(n), rdom(n), tk(0)
    dfn = rev = fa = sdom = dom =
      val = rp = vector<int>(n, -1); }
  void add_edge(int x, int y) { g[x].push_back(y); }
  void dfs(int x) {
    rev[dfn[x] = tk] = x;
    fa[tk] = sdom[tk] = val[tk] = tk; tk++;
    for (int u : g[x]) {
     if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
      r[dfn[u]].push_back(dfn[x]);
    }
  void merge(int x, int y) { fa[x] = y; }
  int find(int x, int c = 0) {
    if (fa[x] == x) return c ? -1 : x;
    if (int p = find(fa[x], 1); p != -1) {
      if (sdom[val[x]] > sdom[val[fa[x]]])
        val[x] = val[fa[x]];
      fa[x] = p;
      return c ? p : val[x];
    } else return c ? fa[x] : val[x];
  vector<int> build(int s) {
    // return the father of each node in dominator tree dfs(s); // p[i] = -2 if i is unreachable, par[s] =
    for (int i = tk - 1; i >= 0; --i) {
      for (int u : r[i])
        sdom[i] = min(sdom[i], sdom[find(u)]);
      if (i) rdom[sdom[i]].push_back(i);
      for (int u : rdom[i]) {
        int p = find(u);
        dom[u] = (sdom[p] == i ? i : p);
      if (i) merge(i, rp[i]);
    vector < int > p(n, -2); p[s] = -1;
    for (int i = 1; i < tk; ++i)</pre>
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < tk; ++i)</pre>
      p[rev[i]] = rev[dom[i]];
    return p;
 }
};
```

5.5 Virtual Tree [f977f8]

```
// copy from 8BQube
vector<int> vG[N];
int top, st[N];
int vrt = -1;
void insert(int u) {
  if (top == -1) return st[++top] = vrt = u, void();
  int p = LCA(st[top], u);
   if(dep[vrt] > dep[p]) vrt = p;
  if (p == st[top]) return st[++top] = u, void();
 while (top >= 1 && dep[st[top - 1]] >= dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
   vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
 top = -1:
  sort(ALL(v),
      [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
 while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do somethina
 reset(vrt);
```

5.6 Fast DMST [7b274d]

```
struct E { int s, t; ll w; }; // 0-base
struct PQ {
  struct P {
    11 v; int i;
    bool operator>(const P &b) const { return v > b.v;
  };
  priority_queue<P, vector<P>, greater<>> pq; 11 tag;
      // min heap
  void push(P p) { p.v -= tag; pq.emplace(p); }
  P top() { P p = pq.top(); p.v += tag; return p; }
  void join(PQ &b) {
    if (pq.size() < b.pq.size())</pre>
      swap(pq, b.pq), swap(tag, b.tag);
    while (!b.pq.empty()) push(b.top()), b.pq.pop();
}; // O(E log^2 V), use leftist tree for O(E log V)
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<PQ> h(n * 2);
  for (int i = 0; i < int(e.size()); ++i)</pre>
   h[e[i].t].push({e[i].w, i});
  vector<int> a(n * 2); iota(iter(a), 0);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
auto o = [&](auto Y, int x) -> int {
    return x==a[x] ? x : a[x] = Y(Y, a[x]); };
  auto S = [&](int i) { return o(o, e[i].s); };
  int pc = v[root] = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
    for (int p = i; v[p]<0 \mid | v[p]==i; p = S(r[p])) {
      if (v[p] == i)
        for (int q = pc++; p != q; p = S(r[p])) {
          h[p].tag -= h[p].top().v; h[q].join(h[p]);
          pa[p] = a[p] = q;
      while (S(h[p].top().i) == p) h[p].pq.pop();
      v[p] = i; r[p] = h[p].top().i;
  vector<int> ans;
  for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
    for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
      v[f] = n;
    ans.push_back(r[i]);
  return ans; // default minimize, returns edgeid array
}
```

5.7 Vizing [58a6ca]

```
// find D+1 edge coloring of a graph with max deg D, O(
             nm)
struct Vizing { // returns maxdeg+1 edge coloring in
              adjacent matrix G
      int n; // 1-based for vertices and colors, simple
                    graph
      vector<vector<int>> C, G;
      vector<int> X, vst;
      Vizing(int _n): n(_n),
      C(n + 1, vector < int > (n + 2)), G(n + 1, vector < int > (n + 1, 
                    + 1)),
      X(n + 1, 1), vst(n + 1) {}
      void solve(vector<pii> &E) {
             auto update = [&](int u)
             { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
             auto color = [&](int u, int v, int c) {
                    int p = G[u][v];
                    G[u][v] = G[v][u] = c;
                    C[u][c] = v, C[v][c] = u;
                    C[u][p] = C[v][p] = 0;
                    if (p) X[u] = X[v] = p;
                    else update(u), update(v);
                    return p;
             auto flip = [&](int u, int c1, int c2) {
                    int p = C[u][c1];
                    swap(C[u][c1], C[u][c2]);
                    if (p) G[u][p] = G[p][u] = c2;
                    if (!C[u][c1]) X[u] = c1;
                    if (!C[u][c2]) X[u] = c2;
                    return p;
              for (int t = 0; t < SZ(E); ++t) {</pre>
```

```
int u = E[t].ff, v0 = E[t].ss, v = v0, c0 = X[u],
          c = c0, d;
     vector<pii> L;
      fill(iter(vst), 0);
     while (!G[u][v0]) {
        L.emplace_back(v, d = X[v]);
        if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
            --a) c = color(u, L[a].ff, c);
        else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
             0; --a) color(u, L[a].ff, L[a].ss);
        else if (vst[d]) break;
        else vst[d] = 1, v = C[u][d];
      if (!G[u][v0]) {
        for (; v; v = flip(v, c, d), swap(c, d));
        if (int a; C[u][c0]) {
          for (a = SZ(L) - 2; a >= 0 && L[a].ss != c;
              --a);
          for (; a >= 0; --a) color(u, L[a].ff, L[a].ss
       else --t;
     }
   }
 }
}:
```

```
5.8 Maximum Clique [d50aa9]
struct MaxClique { // fast when N <= 100</pre>
  bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r), [\&](int x, int y) \{ return d[x] > d[
          y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1:
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < 1ft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
           ]._Find_next(p))
    r[tp] = p, c[tp] = k, ++tp;
dfs(r, c, 1 + 1, mask);
  void dfs(vector<int> &r, vector<int> &c, int 1,
      bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
};
```

5.9 Number of Maximal Clique [11fa26]

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
        if (g[v][some[d][j]])
          some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
        if (g[v][none[d][j]])
          none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
};
```

Minimum Mean Cycle [3e5d2b] 5.10

```
// from 8BQube
11 road[N][N]; // input here
struct MinimumMeanCycle {
  11 dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
for (int i = 2; i <= L; ++i)
      for (int k = 0; k < n; ++k)
         for (int j = 0; j < n; ++j)</pre>
           dp[i][j] =
             min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
      if (dp[L][i] >= INF) continue;
      11 ta = 0, tb = 1;
      for (int j = 1; j < n; ++j)</pre>
         if (dp[j][i] < INF &&</pre>
           ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
           ta = dp[L][i] - dp[j][i], tb = L - j;
      if (ta == 0) continue;
      if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      11 g = \_gcd(a, b);
      return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    for (int i = 0; i < n; ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
  }
};
```

Minimum Steiner Tree [21acea] 5.11

```
// from 8BQube
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
```

```
int vcost[N]; // the cost of vertexs
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
   }
  }
  void add_edge(int ui, int vi, int wi) {
   dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
          dst[i][j] =
            min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)</pre>
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
             submsk = (submsk - 1) \& msk)
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
              vcost[i]);
      for (int i = 0; i < n; ++i) {
        tdst[i] = INF;
        for (int j = 0; j < n; ++j)</pre>
          tdst[i] =
            min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    }
    int ans = INF;
    for (int i = 0; i < n; ++i)</pre>
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
 }
};
5.12 Count Cycles [c7e8f2]
```

```
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
 for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
 for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) c4 += vis[z]++;
  for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M))
```

Math

ll ifloor(ll a,ll b){

6.1 Extended Euclidean Algorithm [c51ae9]

```
// ax+nv = 1, ax+nv == ax == 1 \ (mod \ n)
void extgcd(ll x,ll y,ll &g,ll &a,ll &b) {
 if (y == 0) g=x,a=1,b=0;
  else extgcd(y,x%y,g,b,a),b-=(x/y)*a;
6.2 Floor & Ceil [134881]
```

return a / b - (a % b && (a < 0) ^ (b < 0));

```
ll iceil(ll a,ll b){
 return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre [4e4b23]
// the Jacobi symbol is a generalization of the
    Legendre symbol,
// such that the bottom doesn't need to be prime.
// (n|p) -> same as Legendre
// (n|ab) = (n|a)(n|b)
// work with long long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
}
// 0: a == 0
// -1: a isn't a quad res of p
// else: return X with X^2 % p == a
// doesn't work with long long
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
  int b, d;
  for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
    )) % p;
f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  }
  return g0;
}
6.4 Simplex [aa7741]
 // maximize c^T x
// subject to Ax <= b, x >= 0
// and stores the solution;
typedef long double T; // long double, Rational, double
     + mod<P>...
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-9, inf = 1/.0;
#define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
    s])) s=j
#define rep(i, l, n) for(int i = l; i < n; i++)
struct LPSolver {
  int m, n;
  vector<int> N, B;
  vvd D;
  LPSolver(const vvd& A, const vd& b, const vd& c) :
    m(SZ(b)), n(SZ(c)), N(n+1), B(m), D(m+2, vd(n+2)) {
      rep(i,0,m) \ rep(j,0,n) \ D[i][j] = A[i][j];
      rep(i,0,m) { B[i] = n+i; D[i][n] = -1; D[i][n+1]
          = b[i];}
      rep(j,0,n) \{ N[j] = j; D[m][j] = -c[j]; \}
      N[n] = -1; D[m+1][n] = 1;
```

```
if (s == y) return 100;
  void pivot(int r, int s) {
                                                                     int p = 100 + DiscreteLog(s, x, y, m);
     T *a = D[r].data(), inv = 1 / a[s];
                                                                     if (fpow(x, p, m) != y) return -1;
     rep(i,0,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
                                                                     return p; //returns: x^p = y \pmod{m}
       T *b = D[i].data(), inv2 = b[s] * inv;
       rep(j,0,n+2) b[j] -= a[j] * inv2;
                                                                   6.7 Miller Rabin & Pollard Rho [d3ecd2]
       b[s] = a[s] * inv2;
                                                                   // n < 4,759,123,141
                                                                                                 3 : 2, 7, 61
     rep(j,0,n+2) if (j != s) D[r][j] *= inv;
                                                                   // n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
    rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
                                                                   // n < 2^64
     swap(B[r], N[s]);
                                                                   // 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                   11 mul(ll a, ll b, ll n){
                                                                     return (__int128)a * b % n;
  bool simplex(int phase) {
    int x = m + phase - 1;
                                                                   bool Miller_Rabin(ll a, ll n) {
   if ((a = a % n) == 0) return 1;
     for (;;) {
       int s = -1;
                                                                     if (n % 2 == 0) return n == 2;
       rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
                                                                     ll tmp = (n - 1) / ((n - 1) & (1 - n));
ll t = __lg(((n - 1) & (1 - n))), x = 1;
       if (D[x][s] >= -eps) return true;
       int r = -1;
                                                                     for (; tmp; tmp >>= 1, a = mul(a, a, n))
       rep(i,0,m) {
                                                                       if (tmp \& 1) x = mul(x, a, n);
         if (D[i][s] <= eps) continue;</pre>
                                                                     if (x == 1 || x == n - 1) return 1;
         if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])
                                                                     while (--t)
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
                                                                       if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                     return 0;
       if (r == -1) return false;
       pivot(r, s);
                                                                   bool prime(ll n){
     }
                                                                     vector<ll> tmp = {2, 325, 9375, 28178, 450775,
                                                                          9780504, 1795265022};
                                                                     for(ll i : tmp)
  T solve(vd &x) {
                                                                       if(!Miller_Rabin(i, n)) return false;
     int r = 0;
                                                                     return true;
     rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
     if (D[r][n+1] < -eps) {</pre>
                                                                   map<11, int> cnt;
       pivot(r, n);
                                                                   void PollardRho(ll n) {
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                     if (n == 1) return;
           inf;
                                                                     if (prime(n)) return ++cnt[n], void();
       rep(i,0,m) if (B[i] == -1) {
                                                                     if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
                                                                          void();
         rep(j,1,n+1) ltj(D[i]);
                                                                     11 x = 2, y = 2, d = 1, p = 1;
         pivot(i, s);
                                                                   #define f(x, n, p) ((mul(x, x, n) + p) % n)
       }
                                                                     while (true) {
                                                                       if (d != n && d != 1) {
     bool ok = simplex(1); x = vd(n);
                                                                          PollardRho(n / d);
     rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
                                                                          PollardRho(d);
     return ok ? D[m][n+1] : inf;
                                                                          return;
  }
                                                                       if (d == n) ++p;
                                                                       x = f(x, n, p), y = f(f(y, n, p), n, p);
       Simplex Construction
                                                                       d = gcd(abs(x - y), n);
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for
all 1 \leq j \leq m and x_i \geq 0 for all 1 \leq i \leq n.
1. In case of minimization, let c_i' = -c_i
                                                                   6.8 XOR Basis [006505]
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j \to \mathsf{add} \subseteq \mathsf{and} \supseteq .
                                                                   const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                   struct Basis{
6.6 DiscreteLog [da27bf]
                                                                     int total = 0, rank = 0;
                                                                     vector<ll> b;
int DiscreteLog(int s, int x, int y, int m) {
                                                                     Basis(): b(digit) {}
  constexpr int kStep = 32000;
                                                                     bool add(ll v){ // Gauss Jordan Elimination
  unordered_map<int, int> p;
                                                                       total++:
  int b = 1;
                                                                        for(int i = digit - 1; i >= 0; i--){
  for (int i = 0; i < kStep; ++i) {</pre>
                                                                          if(!(1LL << i & v)) continue;</pre>
    p[y] = i;
y = 1LL * y * x % m;
                                                                          if(b[i] != 0){
                                                                            v ^= b[i];
     b = 1LL * b * x % m;
                                                                            continue;
  for (int i = 0; i < m + 10; i += kStep) {</pre>
                                                                          for(int j = 0; j < i; j++)</pre>
    s = 1LL * s * b % m;
                                                                            if(1LL << j & v) v ^= b[j];</pre>
     if (p.find(s) != p.end()) return i + kStep - p[s];
                                                                          for(int j = i + 1; j < digit; j++)</pre>
                                                                            if(1LL << i & b[j]) b[j] ^= v;</pre>
  return -1;
                                                                          b[i] = v;
                                                                          rank++;
int DiscreteLog(int x, int y, int m) {
                                                                          return true;
  if (m == 1) return 0;
  int s = 1;
                                                                       return false;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
                                                                     11 \text{ getmax}(11 \text{ x} = 0){
     s = 1LL * s * x % m;
```

for(ll i : b) $x = max(x, x ^ i);$

return x;

```
ll getmin(ll x = 0){
    for(ll i : b) x = min(x, x ^ i);
    return x;
  bool can(ll x){
    return getmin(x) == 0;
  11 kth(11 k){ // kth smallest, 0-indexed
    vector<ll> tmp;
    for(11 i : b) if(i) tmp.pb(i);
    11 \text{ ans} = 0;
    for(int i = 0; i < SZ(tmp); i++)</pre>
      if(1LL << i & k) ans ^= tmp[i];</pre>
    return ans;
 }
};
6.9
     Linear Equation [ab487b]
vector<int> RREF(vector<vector<11>> &mat){
  int N = mat.size(), M = mat[0].size();
  int rk = 0;
  vector<int> cols;
  for (int i = 0;i < M;i++) {</pre>
    int cnt = -1;
    for (int j = N-1; j >= rk; j--)
      if(mat[j][i] != 0) cnt = j;
    if(cnt == -1) continue;
    swap(mat[rk], mat[cnt]);
    ll lead = mat[rk][i];
    for (int j = 0; j < M; j++) mat[rk][j] = mat[rk][j] *</pre>
         modinv(lead) % mod;
    for (int j = 0; j < N; j++) {</pre>
      if(j == rk) continue;
      11 tmp = mat[j][i];
      for (int k = 0; k < M; k++)
        mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % mod
              + mod) % mod;
    cols.pb(i);
    rk++;
  return cols;
struct LinearEquation{
  bool ok:
  vector<11> par; //particular solution (Ax = b)
  vector<vector<ll>>> homo; //homogenous (Ax = 0)
  vector<vector<ll>> rref;
  //first M columns are matrix A
  //last column of eq is vector b
  void solve(const vector<vector<ll>> &eq){
    int M = (int)eq[0].size() - 1;
    rref = eq;
    auto piv = RREF(rref);
    int rk = piv.size();
    if(piv.size() && piv.back() == M){
      ok = 0; return;
    ok = 1;
    par.resize(M);
    vector<bool> ispiv(M);
    for (int i = 0;i < rk;i++) {</pre>
      par[piv[i]] = rref[i][M];
      ispiv[piv[i]] = 1;
    for (int i = 0; i < M; i++) {
      if (ispiv[i]) continue;
      vector<ll> h(M);
      h[i] = 1;
      for (int j = 0; j < rk; j++) h[piv[j]] = rref[j][i]</pre>
           ? mod-rref[j][i] : 0;
      homo.pb(h);
 }
};
      Chinese Remainder Theorem [6ef4a3]
```

```
pll solve_crt(ll x1, ll m1, ll x2, ll m2){
 ll g = gcd(m1, m2);
  if ((x2 - x1) % g) return {0, 0}; // no sol
```

```
m1 /= g; m2 /= g;
    11 _, p, q;
    extgcd(m1, m2, _, p, q); // p <= C
ll lcm = m1 * m2 * g;
    ll res = ((__int128)p * (x2 - x1) % lcm * m1 % lcm +
           x1) % lcm;
    // be careful with overflow, C^3
    return {(res + lcm) % lcm, lcm}; // (x, m)
6.11 Sqrt Decomposition [8d7bc0]
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){
   int x = ifloor(n, 1);
    r = ifloor(n, x);
// for all i in [l, r], ceil(n / i) = x
for(int 1, r = n; r >= 1; r = 1 - 1){
   int x = iceil(n, r);
    l = iceil(n, x);
6.12 Floor Sum
• m = |\frac{an+b}{a}|
 • Time complexity: O(\log n)
       f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
                             \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), \\ 0. \end{cases}
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                       \lfloor \tfrac{a}{c} \rfloor \cdot \tfrac{n(n+1)(2n+1)}{6} + \lfloor \tfrac{b}{c} \rfloor \cdot \tfrac{n(n+1)}{2}
                                                                                      a \geq c \vee b \geq c
                                                                                      n < 0 \lor a = 0
                       \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
h(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2
                      \left( \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right)
                      \left[ + \left\lfloor \frac{a}{c} \right\rfloor \cdot \left\lfloor \frac{b}{c} \right\rfloor \cdot \stackrel{\circ}{n}(n+1) \right]
                       +h(a\ \mathsf{mod}\ c,b\ \mathsf{mod}\ c,c,n)
                      +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                       +2\lfloor \frac{\tilde{b}}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                     a \geq c \vee b \geq c
                                                                                      n < 0 \lor a = 0
                       nm(m+1) - 2g(c, c-b-1, a, m-1)
```

Polynomial

7.1 FWHT [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)</pre>
     for (int i = 0; i < n; i += L)</pre>
       for (int j = i; j < i + (L >> 1); ++j)
a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
     1;
void subset_convolution(int *a, int *b, int *c, int L)
     {
  //c_k = \sum_{i=0}^{k} i \cdot j = k, i \cdot i \cdot j = 0
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
```

-2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), otherwise

```
f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
  for (int j = 0; j <= i; ++j)</pre>
                                                                               += P:
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
                                                                        }
  for (int i = 0; i <= L; ++i) fwt(h[i], n,</pre>
                                                                      }
                                                -1):
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
                                                                    if (inv) {
7.2 FFT [13ec2f]
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double
                                                                 }
using val_t = complex<double>;
                                                               };
template<int MAXN>
struct FFT {
  const double PI = acos(-1);
                                                                // Copy from 8BQube
  val_t w[MAXN];
  FFT() {
    for (int i = 0; i < MAXN; ++i) {
  double arg = 2 * PI * i / MAXN;</pre>
                                                                    ++i)
      w[i] = val_t(cos(arg), sin(arg));
                                                               #define V (*this)
    }
  void bitrev(vector<val_t> &a, int n) //same as NTT
  void trans(vector<val_t> &a, int n, bool inv = false)
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
           val_t = a[j + dl] * (inv ? conj(w[x]) : w
               [x]);
           a[j + dl] = a[j] - tmp;
           a[j] += tmp;
                                                                    return V;
      }
                                                                  Poly &imul(ll k) {
    if (inv) {
                                                                    return V;
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
                                                                    int m = 1;
  //multiplying two polynomials A * B:
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
                                                                    assert(m <= MAXN);</pre>
  //A[i] *= B[i], fft.trans(A, siz, 1);
7.3 NTT [bf683f]
                                                                    ntt(X, m, true);
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
                                                                    int m = 1;
// only works when sz(A) + sz(B) - 1 <= MAXN
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
                                                                    assert(m <= MAXN);</pre>
struct NTT {
  11 w[MAXN];
                                                                    Poly Y(V, m);
  11 mpow(ll a, ll n);
  11 minv(ll a) { return mpow(a, P - 2); }
                                                                    fi(0, m) {
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
                                                                    ntt(Xi, m, true);
                                                                    return Xi.isz(n());
  void bitrev(vector<ll> &a, int n) {
    int i = 0;
                                                                    int n = V.n();
    for (int j = 1; j < n - 1; ++j) {
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
                                                                    fc[0] = ifc[0] = 1;
      if (j < i) swap(a[i], a[j]);</pre>
  void operator()(vector<ll> &a, int n, bool inv =
      false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
                                                                    Poly g(n);
      int dx = MAXN / L, dl = L >> 1;
                                                                    11 cp = 1;
      for (int i = 0; i < n; i += L) {</pre>
```

```
for (int j = i, x = 0; j < i + dl; ++j, x += dx
    ll tmp = a[j + dl] * w[x] % P;
    if ((a[j + dl] = a[j] - tmp) < 0) a[j + dl]
    if ((a[j] += tmp) >= P) a[j] -= P;
reverse(a.begin()+1, a.begin()+n);
11 invn = minv(n);
for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
```

7.4 Polynomial Operation [77a8a8]

```
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
#define neg(x) (x ? P - x : 0)
template <int MAXN, 11 P, 11 RT>
                                    // MAXN = 2^k
struct Poly : vector<ll> {
                                  // coefficients in [0,
  using vector<11>::vector;
  static inline NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); }
                                            // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
       .data(), min(p.n(), m), data()); }
  Poly &irev() { return reverse(data(), data() + n()),
  Poly &isz(int m) { return resize(m), V; }
  static ll minv(ll x) { return ntt.minv(x); }
Poly &iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
    fi(0, n()) V[i] = V[i] * k % P;
  Poly Mul(const Poly &rhs) const {
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    Poly X(V, m), Y(rhs, m);
    ntt(X, m), ntt(Y, m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // V[0] != 0, 2*sz<=MAXN
    if (n() == 1) return {minv(V[0])};
    while (m < n() * 2) m <<= 1;</pre>
    Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);
    ntt(Xi, m), ntt(Y, m);
      Xi[i] *= (2 - Xi[i] * Y[i]) % P;
      if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  Poly &shift_inplace(const 11 &c) { // 2 * sz <= MAXN</pre>
    vector<ll> fc(n), ifc(n);
    for (int i = 1; i < n; i++) {
  fc[i] = fc[i - 1] * i % P;</pre>
      ifc[i] = minv(fc[i]);
    for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
```

```
for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
       cp = cp * c % P;
  V = V.irev().Mul(g).isz(n).irev();
  for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
      Ρ;
  return V;
Poly shift(const ll &c) const { return Poly(V).
    shift_inplace(c); }
Poly _Sqrt() const { // Jacobi(V[0], P) = 1
  if (n() == 1) return {QuadraticResidue(V[0], P)};
  Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
      1);
Poly Sqrt() const { // 2 * sz <= MAXN
  Poly a;
  bool has = 0;
  for (int i = 0; i < n(); i++) {</pre>
    if (V[i]) has = 1;
    if (has) a.push_back(V[i]);
  if (!has) return V;
  if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
    return Poly();
  a = a.isz((n() + a.n()) / 2)._Sqrt();
  int sz = a.n();
  a.isz(n());
  rotate(a.begin(), a.begin() + sz, a.end());
pair<Poly, Poly> DivMod(const Poly &rhs) const { //
    (rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, V};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs);
  X.irev().isz(m);
  Poly Y(V);
  Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
 X = rhs.Mul(Q), Y = V;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;</pre>
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
  return ret;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const vector<ll> &x, const vector<
    Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
      .second:
  down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
       _tmul(m, V);
  fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
      1, down[i / 2]);
  vector<11> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = {neg(x[i]), 1};
  for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
      Mul(up[i * 2 + 1]);
  return up;
```

```
vector<ll> Eval(const vector<ll> &x) const { // 1e5,
     15
  auto up = _tree1(x);
 return _eval(x, up);
static Poly Interpolate(const vector<11> &x, const
    vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
 vector<Poly> up = _tree1(x), down(m * 2);
vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m - 1; i > 0; --i)
    down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
        down[i * 2 + 1].Mul(up[i * 2]));
 return down[1];
Poly Ln() const { //V[0] == 1, 2*sz <= MAXN
 return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { //V[0] == 0,2*sz <= MAXN
  if (n() == 1) return {1};
  Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln();
  Y[0] = P - 1;
  fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
  return X.Mul(Y).isz(n());
// M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(11 k) const { // 2*sz<=MAXN
  int nz = 0;
  while (nz < n() && !V[nz]) ++nz;</pre>
  if (nz * min(k, (11)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly{1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const ll c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
      n()).irev();
// sum_j w_j [x^j] f(x^i) for i \in [0, m]
Poly power_projection(Poly wt, int m) { // 4*sz <=
    MAXN!
  assert(n() == wt.n());
  if (!n()) {
    return Poly(m + 1, 0);
  if (V[0] != 0) {
    11 c = V[0];
    V[0] = 0;
    Poly A = V.power_projection(wt, m);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
        factorial
    Poly B(m + 1);
    11 pow = 1;
    fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
        * c % P; // inv. of fac
    A = A.Mul(B).isz(m + 1);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P;
    return A;
  int n = 1;
  while (n < V.n()) n *= 2;</pre>
  isz(n), wt.isz(n).irev();
  int k = 1;
  Poly p(wt, 2 * n), q(V, 2 * n);
  q.imul(P - 1);
 while (n > 1) {
   Poly r(2 * n * k);
    fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
        q[i]));
    Poly pq = p.Mul(r).isz(4 * n * k);
    Poly qq = q.Mul(r).isz(4 * n * k);
    fi(0, 2 * n * k) {
      pq[2 * n * k + i] += p[i];
      qq[2 * n * k + i] += q[i] + r[i];
      pq[2 * n * k + i] %= P;
      qq[2 * n * k + i] \% = P;
    fill(p.begin(), p.end(), 0);
```

```
fill(q.begin(), q.end(), 0);
      for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
  p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];
        q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];
      n /= 2, k *= 2;
    }
    Poly ans(k);
    fi(0, k) ans[i] = p[2 * i];
    return ans.irev().isz(m + 1);
  Poly FPSinv() {
    const int n = V.n() - 1;
    if (n == -1) return {};
    assert(V[0] == 0);
    if (n == 0) return V;
    assert(V[1] != 0);
    ll c = V[1], ic = minv(c);
    imul(ic);
    Poly wt(n + 1);
    wt[n] = 1;
    Poly A = V.power_projection(wt, n);
    Poly g(n);
    fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
    g = g.Pow(neg(minv(n)));
    g.insert(g.begin(), 0);
    11 pow = 1;
    fi(0, g.n()) g[i] = g[i] * pow % P, pow = pow * ic
        % P;
    return g;
  Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
      1 = k;
    return Poly(*this).irev().Mul(rhs).isz(n()).irev();
  Poly FPScomp(Poly g) { // solves V(g(x))
    auto rec = [&](auto &rec, int n, int k, Poly Q) ->
        Poly {
      if (n == 1) {
        Poly p(2 * k);
        irev();
        fi(0, k) p[2 * i] = V[i];
        return p;
      Poly R(2 * n * k);
      fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
           Q[i]));
      Poly QQ = Q.Mul(R).isz(4 * n * k);
fi(0, 2 * n * k) {
        QQ[2 * n * k + i] += Q[i] + R[i];
        QQ[2 * n * k + i] %= P;
      Poly nxt_Q(2 * n * k);
for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
         nxt_{Q[n * j + i]} = QQ[(2 * n) * j + (2 * i + 0)
      Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
      Poly pq(4 * n * k);

for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {

pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +
             i];
        pq[(2*n)*j+(2*i+1)]%= P;
      Poly p(2 * n * k);
      fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i
           1) % P;
      pq.pop_back();
      Poly x = pq.TMul(R);
fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
      return p;
    };
    int sz = 1;
    while(sz < n() || sz < g.n()) sz <<= 1;
    return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
          sz)).isz(sz).irev();
 }
};
#undef fi
```

#undef V

```
#undef neg
using Poly_t = Poly<1 << 19, 998244353, 3>;
```

7.5 Generating Function **Ordinary Generating Function**

- C(x) = A(rx): $c_n = r^n a_n$ 的一般生成函數。
- C(x) = A(x) + B(x): $c_n = a_n + b_n$ 的一般生成函數。
- C(x) = A(x)B(x): $c_n = \sum_{i=0}^n a_i b_{n-i}$ 的一般生成函數。
 $C(x) = A(x)^k$: $c_n = \sum_{i_1+i_2+\dots+i_k=n} a_{i_1}a_{i_2}\dots a_{i_k}$ 的一 $a_{i_1}a_{i_2}\dots a_{i_k}$ 的一般生成函數。
- C(x) = xA(x)': $c_n = na_n$ 的一般生成函數。
- $C(x) = \frac{A(x)}{1-x}$: $c_n = \sum_{i=0}^n a_i$ 的一般生成函數。
- $C(x)=A(1)+x\frac{A(1)-A(x)}{1-x}$: $c_n=\sum\limits_{i=n}^{\infty}a_i$ 的一般生成函數。

常用展開式

- $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$, ${a \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

• 卡特蘭數: $f(x) = \frac{1 - \sqrt{1 - 4x}}{2}$

Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數 $\hat{C}(x) = \hat{A}(x)^k$: $\sum_{i_1+i_2+\dots+i_k=n} \binom{n}{(i_1,i_2,\dots,i_k)} a_i a_{i_2} \dots a_{i_k}$ 的指數生成函數
- $\hat{C}(x) = \exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 $F(0), G(0) = 0, F'(0), G'(0) \neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$
$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

7.6 Bostan Mori [41c3bc]

```
const 11 mod = 998244353;
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = \{P, Q\};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  vector<11> Qn(sz);
  vector N(2, vector<ll>(sz));
  while(k) {
    fill(iter(Qn), 0);
for(int i = 0; i < d + 1; i++){</pre>
      Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    ntt(Qn, sz, false);
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){</pre>
      fill(iter(N[i]), 0);
      copy(iter(M[i]), N[i].begin());
      ntt(N[i], sz, false);
      for(int j = 0; j < sz; j++)</pre>
        N[i][j] = N[i][j] * Qn[j] % mod;
      ntt(N[i], sz, true);
      for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
        M[i][j >> 1] = N[i][j];
```

```
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      }
    }
    k >>= 1;
  return M[0][0] * ntt.minv(M[1][0]) % mod;
11 LinearRecursion(vector<11> a, vector<11> c, 11 k) {
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (lig(2 * d) + lig(2 * d)))
      1)));
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
  auto A = a; A.resize(sz);
  auto C = c; C.resize(sz);
  ntt(A, sz, false), ntt(C, sz, false);
  for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
  ntt(A, sz, true);
  A.resize(d);
  return BostanMori(A, c, k);
}
8
    String
8.1 KMP Algorithm [c8b75f]
// 0-based
// fail[i] = max k<i s.t. s[0..k] = s[i-k..i]
vector<int> kmp_build_fail(const string &s){
  int n = SZ(s);
  vector<int> fail(n, -1);
  int cur = -1;
  for(int i = 1; i < n; i++){</pre>
    while(cur != -1 && s[cur + 1] != s[i])
```

```
cur = fail[cur];
    if(s[cur + 1] == s[i])
      cur++;
    fail[i] = cur;
  }
  return fail;
void kmp_match(const string &s, const vector<int> &fail
    , const string &t){
  int cur = -1;
  int n = SZ(s), m = SZ(t);
  for(int i = 0; i < m; i++){</pre>
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t
        [i]))
      cur = fail[cur];
    if(cur + 1 < n && s[cur + 1] == t[i])
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
}
```

8.2 Manacher Algorithm [caf0f4]

8.3 Lyndon Factorization [7c612b]

```
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
```

```
// each w[i] strictly smaller than all its suffix
void duval(const string &s, vector<pii> &w) {
  for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
    for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
k = (s[k] < s[j] ? i : k + 1);</pre>
     // if (i < n / 2 && j >= n / 2) {
     // for min cyclic shift, call duval(s + s)
// then here s.substr(i, n / 2) is min cyclic shift
     for (; i <= k; i += j - k)</pre>
       w.pb(pii(i, j - k)); // s.substr(l, len)
  }
}
8.4 Suffix Array [cd67ea]
struct SuffixArray {
  vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]
       and sa[i-1]
                                  // sa[0] = s.size()
                                  // character should be 1-
                                       based
  SuffixArray(string& s, int lim=256) { // or
       basic_string<int>
     int n = s.size() + 1, k = 0, a, b;
     vector<int> x(n, 0), y(n), ws(max(n, lim));
     rank.assign(n, 0);
     for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
    sa = lcp = y, iota(sa.begin(), sa.end(), 0);
for (int j = 0, p = 0; p < n; j = max(1, j * 2),</pre>
         lim = p) {
       p = j, iota(y.begin(), y.end(), n - j);
       for (int i = 0; i < n; i++)
  if (sa[i] >= j) y[p++] = sa[i] - j;
       for (int &i : ws) i = 0;
       for (int i = 0; i < n; i++) ws[x[i]]++;
for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];</pre>
       for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
       swap(x, y), p = 1, x[sa[0]] = 0;
       for(int i = 1; i < n; i++){</pre>
         a = sa[i - 1], b = sa[i];
         x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ?
               p - 1 : p++;
       }
     for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
     for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
       for (k && k--, j = sa[rank[i] - 1];
            s[i + k] == s[j + k]; k++);
  }
};
8.5
     Suffix Automaton [016373]
struct exSAM {
  const int CNUM = 26;
  // len: maxlength, link: fail link
  // LenSorted: topo order, cnt: occur
  vector<int> len, link, lenSorted, cnt;
  vector<vector<int>> next;
  int total = 0;
  int newnode() {
     return total++;
  void init(int n) { // total number of characters
  len.assign(2 * n, 0); link.assign(2 * n, 0);
     lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
     next.assign(2 * n, vector<int>(CNUM));
    newnode(), link[0] = -1;
  int insertSAM(int last, int c) {
```

// not exSAM: cur = newnode(), p = last

next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;

if (len[p] + 1 == len[q]) return link[cur] = q, cur

int cur = next[last][c];

int p = link[last];

int q = next[p][c];

int clone = newnode();

len[cur] = len[last] + 1;

while (p != -1 && !next[p][c])

for (int i = 0; i < CNUM; ++i)</pre>

```
next[clone][i] = len[next[q][i]] ? next[q][i] :
                                                                 else rep.emplace_back(Rep({sft + c - R - l + 1, sft
                                                                      + c - L - 1 + 1, 1));
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
                                                               for (int cntr = 0; cntr < n; cntr++) {</pre>
      next[p][c] = clone, p = link[p];
                                                                 int 1, k1, k2;
    link[link[cur] = clone] = link[q];
                                                                 if (cntr < nu) {</pre>
    link[q] = clone;
                                                                   1 = nu - cntr;
    return cur;
                                                                   k1 = get_z(z1, nu - cntr);
                                                                   k2 = get_z(z2, nv + 1 + cntr);
  void insert(const string &s) {
                                                                 } else {
    int cur = 0;
                                                                   l = cntr - nu + 1;
                                                                   k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
                                                                   k2 = get_z(z4, (cntr - nu) + 1);
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
                                                                 if (k1 + k2 >= 1)
                                                                   add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
    }
                                                               }
  void build() {
                                                             }
    queue<int> q;
                                                             8.8 AC Automaton [f529e6]
    q.push(0);
    while (!q.empty()) {
                                                             const int SIGMA = 26;
      int cur = q.front();
                                                             struct AC_Automaton {
      q.pop();
                                                               // child: trie, next: automaton
      for (int i = 0; i < CNUM; ++i)</pre>
                                                               vector<vector<int>> child, next;
        if (next[cur][i])
                                                               vector<int> fail, cnt, ord;
          q.push(insertSAM(cur, i));
                                                               int total = 0;
    }
                                                               int newnode() {
    vector<int> lc(total);
                                                                 return total++;
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
                                                               void init(int len) { // len >= 1 + total len
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
                                                                 child.assign(len, vector<int>(26, -1));
        i]]] = i;
                                                                 next.assign(len, vector<int>(26, -1));
                                                                 fail.assign(len, -1); cnt.assign(len, 0);
  void solve() {
                                                                 ord.clear();
    for (int i = total - 2; i >= 0; --i)
                                                                 newnode();
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
 }
                                                               int input(string &s) {
};
                                                                 int cur = 0;
8.6 Z-value Algorithm [488d87]
                                                                 for (char c : s) {
                                                                   if (child[cur][c - 'A'] == -1)
  child[cur][c - 'A'] = newnode();
cur = child[cur][c - 'A'];
// z[i] = max k s.t. s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z_function(const string &s){
                                                                 return cur; // return the end node of string
 int n = s.size();
  vector<int> z(n);
                                                               void make_fl() {
  for(int i = 1, l = 0, r = 0; i < n; i++){
                                                                 queue<int> q;
    if(i <= r) z[i] = min(r - i + 1, z[i - 1]);</pre>
                                                                 q.push(0), fail[0] = -1;
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
                                                                 while(!q.empty()) {
     z[i]++;
                                                                   int R = q.front();
    if(i + z[i] - 1 > r)
                                                                   q.pop(); ord.pb(R);
for (int i = 0; i < SIGMA; i++)</pre>
      l = i, r = i + z[i] - 1;
 }
                                                                     if (child[R][i] != -1) {
  return z;
                                                                       int X = next[R][i] = child[R][i], Z = fail[R
8.7 Main Lorentz [fcfb8f]
                                                                       while (Z != -1 && child[Z][i] == -1)
                                                                         Z = fail[Z];
struct Rep{ int minl, maxl, len; };
                                                                       fail[X] = Z != -1 ? child[Z][i] : 0;
vector<Rep> rep; // 0-base
                                                                       q.push(X);
// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)
void main_lorentz(const string &s, int sft = 0) {
                                                                     else next[R][i] = R ? next[fail[R]][i] : 0;
  const int n = s.size();
                                                                 }
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
                                                               void solve() {
  for (int i : ord | views::reverse)
  const string u = s.substr(0, nu), v = s.substr(nu),
        ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
                                                                   if (i) cnt[fail[i]] += cnt[i];
            ());
 main_lorentz(u, sft), main_lorentz(v, sft + nu);
                                                             };
  const auto z1 = z_function(ru), z2 = z_function(v + '
      #' + u),
                                                                   Palindrome Automaton [8a071b]
                                                             8.9
             z3 = z function(ru + '#' + rv), z4 =
                 z_function(v);
                                                             struct PalindromicTree {
  auto get_z = [](const vector<int> &z, int i) {
                                                               struct node {
                                                                 int nxt[26], fail, len; // num = depth of fail link
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
                                                                 int cnt, num; // cnt = occur, num = #pal_suffix of
  auto add_rep = [&](bool left, int c, int l, int k1,
                                                                     this node
     int k2) {
                                                                 node(int l = 0) : nxt{},fail(0),len(1),cnt(0),num
    const int L = max(1, 1 - k2), R = min(1 - left, k1)
                                                                     (0) {}
    if (L > R) return;
                                                               vector<node> st; vector<int> s; int last, n;
    if (left) rep.emplace_back(Rep({sft + c - R, sft +
                                                               void init() {
        c - L, 1}));
                                                                 st.clear(); s.clear(); last = 1; n = 0;
```

return S_best;

```
st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
  int getFail(int x) {
   while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x;
  void add(int c) {
    s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
     int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
      st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
 }
  void dpcnt() {
    for(int i = SZ(st) - 1; i >= 0; i--){
     auto nd = st[i];
      st[nd.fail].cnt += nd.cnt;
  int size() { return (int)st.size() - 2; }
};
```

9 Misc

9.1 Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (l + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

9.2 Matroid

 $M=(E,\mathcal{I})$, where $\mathcal{I}\subseteq 2^E$ is nonempty, is a matroid if:

- If $S \in \mathcal{I}$ and $S' \subsetneq S$, then $S' \in \mathcal{I}$.
- For $S_1,S_2\in\mathcal{I}$ s.t. $|S_1|<|S_2|$, there exists $e\in S_2\setminus S_1$ s.t. $S_1\cup\{e\}\in\mathcal{I}$. Matroid intersection:

Start from $S=\emptyset$. In each iteration, let

- $Y_1 = \{x \notin S \mid S \cup \{x\} \in \mathcal{I}_1\}$
- $Y_2 = \{x \notin S \mid S \cup \{x\} \in \mathcal{I}_2\}$

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \not \in S$, create edges

- $x \to y \text{ if } S \{x\} \cup \{y\} \in \mathcal{I}_1.$
- $y \to x$ if $S \{x\} \cup \{y\} \in \mathcal{I}_2$.

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.3 Simulate Annealing [ff826c]

```
Id anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<ld> rnd(0, 1);
  const ld dT = 0.001;
  // Argument p
  ld S_cur = calc(p), S_best = S_cur;
  for (ld T = 2000; T > eps; T -= dT) {
      // Modify p to p_prime
      const ld S_prime = calc(p_prime);
      const ld delta_c = S_prime - S_cur;
      ld prob = min((ld)1, exp(-delta_c / T));
      if (rnd(rnd_engine) <= prob)
            S_cur = S_prime, p = p_prime;
      if (S_prime < S_best) // find min
            S_best = S_prime, p_best = p_prime;
}</pre>
```

```
9.4 Binary Search On Fraction [f6b9ec]
```

```
struct Q {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p * d, q + b.q * d};
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N, auto &&pred) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step /= 2 : step *=</pre>
        2);)
      if (Q mid = hi.go(lo, len + step);
          mid.p > N \mid\mid mid.q > N \mid\mid dir ^ pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

9.5 Min Plus Convolution [09b5c3]

10 Notes

10.1 Geometry

 $a \qquad b \qquad c \qquad 2$ Law of cosines: $a^2 = b^2 + c^2 - 2bc\cos\alpha$

a + b

a - b

Law of tangents:

Rotation Matrix

```
• rotate 90^\circ: (x,y) \to (-y,x)

• rotate -90^\circ: (x,y) \to (y,-x)

Triangles
Side lengths: a,b,c
Semiperimeter: p = \frac{a+b+c}{2}
Area: A = \sqrt{p(p-a)(p-b)(p-c)}
Circumradius: R = \frac{abc}{4A}
Inradius: r = \frac{A}{p}
Length of median (divides triangle into two equal-area triangles): m_a = \frac{1}{2}\sqrt{2b^2+2c^2-a^2}
Length of bisector (divides angles in two): s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}
Law of sines: \frac{\sin\alpha}{2} = \frac{\sin\beta}{2} = \frac{\sin\gamma}{2} = \frac{1}{2}
```

 $\cos \theta$

 $\int \sin \theta$

 $\tan \frac{\alpha + \beta}{\alpha}$

tan -

 $\alpha - \beta$

 $-\sin\theta$

 $\cos \theta$

Ouadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A=2ef\cdot\sin\theta=F\tan\theta=\sqrt{4e^2f^2-F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \mathrm{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \mathrm{atan2}(y,x) \end{array}$$

Green's Theorem

$$\iint_{D} \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy = \oint_{L^{+}} (Pdx + Qdy)$$

$$\mathsf{Area} = \frac{1}{2} \oint_{L} x \, dy - y \, dx$$

Circular sector:

$$\begin{aligned} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r(r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{aligned}$$

Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1,p_2,p_3 are collinear $\iff p_1^*,p_2^*,p_3^*$ intersect at a point p lies above $l \iff l^*$ lies above p^*
- lower convex hull \leftrightarrow upper envelope

10.2 Trigonometry

$$\begin{split} \sinh x &= \frac{1}{2}(e^x - e^{-x}) & \cosh x &= \frac{1}{2}(e^x + e^{-x}) \\ & \sin n\pi &= 0 & \cos n\pi &= (-1)^n \\ & \sin(\alpha + \beta) &= \sin \alpha \cos \beta + \cos \alpha \sin \beta \end{split}$$

$$\begin{aligned} \cos(\alpha+\beta) &= \cos\alpha\cos\beta - \sin\alpha\sin\beta\\ \sin(2\alpha) &= 2\cos\alpha\sin\alpha\\ \cos(2\alpha) &= \cos^2\alpha - \sin^2\alpha\\ &= 2\cos^2\alpha - 1\\ &= 1 - 2\sin^2\alpha \end{aligned}$$

$$\begin{split} \tan(\alpha+\beta) &= \frac{\tan\alpha + \tan\beta}{1-\tan\alpha\tan\beta} \\ \sin\alpha + \sin\beta &= 2\sin\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \cos\alpha + \cos\beta &= 2\cos\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \sin\alpha\sin\beta &= \frac{1}{2}(\cos(\alpha-\beta) - \cos(\alpha+\beta)) \\ \sin\alpha\cos\beta &= \frac{1}{2}(\sin(\alpha+\beta) + \sin(\alpha-\beta)) \\ \cos\alpha\sin\beta &= \frac{1}{2}(\sin(\alpha+\beta) - \sin(\alpha-\beta)) \\ \cos\alpha\cos\beta &= \frac{1}{2}(\cos(\alpha-\beta) + \cos(\alpha+\beta)) \end{split}$$

where V, W are lengths of sides opposite angles α, β .

$$a\cos x + b\sin x = r\cos(x - \phi)$$

 $a\sin x + b\cos x = r\sin(x + \phi)$

 $(V+W)\tan(\alpha-\beta)/2=(V-W)\tan(\alpha+\beta)/2$

where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

10.3 Calculus

Integration by parts:

$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx} \tan x = 1 + \tan^2 x \qquad \frac{d}{dx} \arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int xe^{ax} = \frac{e^{ax}}{a^2} (ax - 1)$$

$$\int \sin^2(x) = \frac{x}{2} - \frac{1}{4} \sin 2x \qquad \int \sin^3 x = \frac{1}{12} \cos 3x - \frac{3}{4} \cos x$$

$$\int \cos^2(x) = \frac{x}{2} + \frac{1}{4} \sin 2x \qquad \int \cos^3 x = \frac{1}{12} \sin 3x + \frac{3}{4} \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int xe^x = e^x(x - 1) \qquad \int x^2 e^x = e^x(x^2 - 2x + 2)$$

$$\int x^2 \sin x = 2x \sin x - (x^2 - 2) \cos x$$

$$\int x^2 \cos x = 2x \cos x + (x^2 - 2) \sin x$$

$$\int e^x \sin x = \frac{1}{2} e^x (\sin x - \cos x)$$

$$\int e^x \cos x = \frac{1}{2} e^x (\sin x + \cos x)$$

$$\int x e^x \sin x = \frac{1}{2} e^x (x \sin x - x \cos x + \cos x)$$

$$\int x e^x \cos x = \frac{1}{2} e^x (x \sin x + x \cos x - \sin x)$$

10.4 Sum & Series

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$\begin{split} e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \, (-\infty < x < \infty) \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \, (-1 < x \le 1) \\ \sqrt{1+x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \, (-1 \le x \le 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \, (-\infty < x < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \, (-\infty < x < \infty) \end{split}$$

10.5 Misc

• Cramer's rule

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.

- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each $\emph{labeled}$ vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \dots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

• Erdős–Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only

if
$$d_1+\cdots+d_n$$
 is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds

for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n

is bigraphic if and only if
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

holds for every $1 \le k \le n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$

$$\cdots \geq a_n \text{ is digraphic if and only if } \sum_{i=1}^n a_i = \sum_{i=1}^n b_i \text{ and } \sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-\frac{n}{2}) \frac{n}{2} = \frac{n}{2} \sum_{i=1}^n a_i = \sum_{i=1}^n a_i$$

$$1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have

 $A=\#\{\mbox{lattice points in the interior}\}+\frac{\#\{\mbox{lattice points on the boundary}\}}{2}-1.$

Möbius inversion formula

-
$$f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$$

-
$$f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$$

- A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap, θ : arcsin(a/r).
- Volume = $\pi h^2 (3r h)/3 = \pi h(3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \sin \theta)$ $\cos \theta$)²/3.
- Area $= 2\pi r h = \pi (a^2 + h^2) = 2\pi r^2 (1 \cos \theta).$

Lagrange multiplier

- Optimize $f(x_1, \ldots, x_n)$ when k constraints $g_i(x_1, \ldots, x_n) = 0$.
- Lagrangian function

 $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n).$

The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.

· Nearest points of two skew lines

- Line 1 : ${m v}_1 = {m p}_1 + t_1 {m d}_1$
- Line 2 : ${m v}_2 = {m p}_2 + t_2 {m d}_2$
- $n = d_1 \times d_2$
- $n_1 = d_1 \times n$

- $egin{aligned} & m{n}_1 = m{a}_1 & \cdots & \ & m{n}_2 = m{d}_2 imes m{n} \ & m{c}_1 = m{p}_1 + rac{(m{p}_2 m{p}_1) \cdot m{n}_2}{m{d}_1 \cdot m{n}_2} m{d}_1 \ & m{c}_2 = m{p}_2 + rac{(m{p}_1 m{p}_2) \cdot m{n}_1}{m{d}_2 \cdot m{n}_1} m{d}_2 \end{aligned}$

• Bernoulli numbers
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, \text{ EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^{\infty} B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{array}{l} S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n = \sum_{i=0}^n S(n,i)(x)_i \\ \bullet \text{ Pentagonal number theorem} \end{array}$$

$$\prod_{n=1}^{\infty} (1 - x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers
$$1$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {n \choose n}$$

$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

• Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$. $E(n, \overline{k}) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$ E(n,0) = E(n,n-1) = 1 $E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$

10.6 Number

· Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039, 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

Number of paritions of n:

 $n \mid 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 20 \ 30 \ 40 \ 50 \ 100$ p(n) 2 3 5 7 11 15 22 30 627 5604 4e4 2e5 2e8

Maximum number of divisors:

n | 100 1e3 1e6 1e9 1e12 1e15 1e18 |d(i)| 12 32 240 1344 6720 26880 103680

n 1234567 8 11 12 13 14 15 $\binom{2n}{n}$ 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

• Number of ways to partition a set of n labeled elements: $n \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \mid 10 \mid 11 \mid 12 \mid 13$ B_n 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7

