Contents 8 Polynomial 19			
	Basic 1.1 Default Code 1.2 .vimrc 1.3 Fast IO 1.4 Random 1.5 Checker 1.6 PBDS Tree Data Structure	1 1 2 2 2 2 2 2	8.1 FWHT 19 8.2 FFT 19 8.3 NTT 15 8.4 Polynomial Operation 20 8.5 Generating Function 21 8.5.1 Ordinary Generating Function 21 8.5.2 Exponential Generating Function 21 8.6 Bostan Mori 21 9 String 22
2	2.1 Heavy-Light Decomposition 2.2 Link Cut Tree 2.3 Treap 2.4 KD Tree 2.5 Leftist Tree	2 2 3 4 4	9.1 KMP Algorithm
	Flow & Matching 3.1 Dinic 3.2 Bounded Flow 3.3 MCMF 3.4 Min Cost Circulation 3.5 Gomory Hu 3.6 Stoer Wagner Algorithm 3.7 Bipartite Matching 3.8 Kuhn Munkres Algorithm 3.9 Max Simple Graph Matching 3.10 Stable Marriage 3.11 Flow Model	4 5 5 6 6 6 6 7 7	9.7 Main Lorentz 23 9.8 AC Automaton 25 9.9 Palindrome Automaton 24 10 Formula 24 10.1 Recurrences 24 10.2 Geometry 24 10.2.1 Rotation Matrix 24 10.2.2 Triangles 24 10.2.3 Quadrilaterals 24 10.2.4 Spherical coordinates 24 10.2.5 Green's Theorem 24 10.2.6 Point-Line Duality 24
4	4.1 Geometry Template 4.2 Convex Hull 4.3 Minimum Enclosing Circle 4.4 Minkowski Sum 4.5 Polar Angle Comparator 4.6 Half Plane Intersection 4.7 Dynamic Convex Hull 4.8 3D Point 4.9 ConvexHull3D 4.10 Circle Operations 4.11 Delaunay Triangulation 4.12 Voronoi Diagram 4.13 Polygon Union 4.14 Tangent Point to Convex Hull	8 8 8 9 9 9 10 10 11 11 12 12 12	10.3 Trigonometry
5	Graph 5.1 Block Cut Tree 5.2 2-SAT 5.3 Dominator Tree 5.4 Virtual Tree 5.5 Directed Minimum Spanning Tree 5.6 Fast DMST 5.7 Vizing 5.8 Maximum Clique 5.9 Number of Maximal Clique 5.10 Minimum Mean Cycle 5.11 Minimum Steiner Tree	12 12 13 13 13 14 14 14 15 15 15	<pre>#define io ios_base::sync_with_stdio(0);cin.tie(0);cerr.tie</pre>
6	Math 6.1 Extended Euclidean Algorithm 6.2 Floor & Ceil 6.3 Legendre 6.4 Simplex 6.5 Floor Sum 6.6 DiscreteLog 6.7 Miller Rabin & Pollard Rho 6.8 XOR Basis 6.9 Linear Equation 6.10 Chinese Remainder Theorem 6.11 Sqrt Decomposition	16 16 16 17 17 17 17 18 18 18	<pre>#ifdef zisk void debug(){cerr << "\n";} template<class class="" t,="" u=""> void debug(T a, U b){cerr << a << " ", debug(b);} template<class t=""> void pary(T l, T r){ while (1 != r) cerr << *l << " ", l++; cerr << "\n"; } #else #define debug() void() #define pary() void() #endif</class></class></pre>
7	Misc 7.1 Cyclic Ternary Search	19 19 19	<pre>template < class A, class B> ostream& operator < < (ostream& o, pair < A, B> p) { return o << '(' << p.ff << ',' << p.ss << ')'; } int main(){ io; }</pre>

1.2 .vimrc

1.3 Fast IO

```
// from JAW
inline int my_getchar() {
  const int N = 1 << 20;
  static char buf[N];
  static char *p = buf , *end = buf;
  if(p == end) {
    if((end = buf + fread(buf , 1 , N , stdin)) == buf)
        return EOF;
    p = buf;
  }
  return *p++;
}
inline int readint(int &x) {
  static char c , neg;
  while((c = my_getchar()) < '-') {</pre>
    if(c == EOF) return 0;
  neg = (c == '-') ? -1 : 1;
  x = (neg == 1) ? c - '0' : 0;
  while((c = my_getchar()) >= '0') x = (x << 3) + (x << 1)
      + (c - '0');
  x *= neg;
  return 1;
const int kBufSize = 524288;
char inbuf[kBufSize];
char buf_[kBufSize]; size_t size_;
inline void Flush_() { write(1, buf_, size_); size_ = 0; }
inline void CheckFlush_(size_t sz) { if (sz + size_ >
    kBufSize) Flush_(); }
inline void PutInt(int a) {
  static char tmp[22] = "01234567890123456789\n";
  CheckFlush_(10);
  if(a < 0){
    *(buf_ + size_) = '-';
    a = ~a + 1;
    size_++;
  int tail = 20;
  if (!a) {
    tmp[--tail] = '0';
  } else {
    for (; a; a /= 10) tmp[--tail] = (a % 10) ^ '0';
  memcpy(buf_ + size_, tmp + tail, 21 - tail);
  size_ += 21 - tail;
int main(){
  Flush_();
  return 0;
```

1.4 Random

1.5 Checker

```
#!/usr/bin/env bash
set -e
while :; do
    python3 gen.py > test.txt
    diff <(./a.exe < test.txt) <(./b.exe < test.txt)
done</pre>
```

1.6 PBDS Tree

2 Data Structure

2.1 Heavy-Light Decomposition

```
struct Heavy_light_Decomposition { // 1-base
  int n, up[maxn], dep[maxn], to[maxn], siz[maxn], pa[maxn
  int C, ti[maxn], ord[maxn], wdown[maxn], edge[maxn], et =
       0;
  vector<pii> G[maxn];
  void init(int _n) {
    n = n, C = 0, et = 1;
    for (int i = 1;i <= n;i++)</pre>
      G[i].clear(), to[i] = 0;
  void add_edge(int a, int b, int w) {
    G[a].push_back(pii(b, et)), G[b].push_back(pii(a, et));
    edge[et++] = w;
  void dfs(int u, int f, int d) {
    siz[u] = 1, pa[u] = f, dep[u] = d;
    for (auto &v: G[u])
      if (v.ff != f)
        dfs(v.ff, u, d+1), siz[u] += siz[v];
        if (siz[to[u]] < siz[v]) to[u] = v;
  void cut(int u, int link) {
    ti[u] = C;
    ord[C++] = u, up[u] = link;
    if (!to[u]) return;
    cut(to[u], link);
    for (auto v:G[u]) {
      if (v.ff != pa[u] && v.ff != to[u]) cut(v.ff, v.ff);
  void build() { dfs(1, 1, 1), cut(1, 1); }
  int query(int a, int b) {
    int ta = up[a], tb = up[b], re = 0;
    while (ta != tb)
      if (dep[ta] < dep[tb])</pre>
        /*query*/, tb = up[b = pa[tb]];
      else /*query*/, ta = up[a = pa[ta]];
    if (a == b) return re;
    if (ti[a] > ti[b]) swap(a, b);
    /*query*/
    return re;
  }
};
```

2.2 Link Cut Tree

```
struct Splay { // LCT + PATH add
    static Splay nil;
```

```
Splay *ch[2], *f;
int rev;
int sz;
11 val, sum, tag;
Splay() : rev(0), sz(1), val(1), sum(1), tag(0) {
  f = ch[0] = ch[1] = &nil;
bool isr() { return f->ch[0] != this && f->ch[1] != this;
int dir() { return f->ch[0] == this ? 0 : 1; }
void setCh(Splay *c, int d) {
  ch[d] = c;
  if (c != &nil) c->f = this;
  pull();
void push() {
  for(int i = 0; i < 2; i++){
    if(ch[i] == &nil) continue;
    if(rev) swap(ch[i]->ch[0], ch[i]->ch[1]), ch[i]->rev
        ^= 1;
    if(tag != 0){
      ch[i]->tag += tag;
      ch[i]->val += tag;
      ch[i]->sum += tag * ch[i]->sz;
    }
  }
  tag = 0;
  rev = 0;
void pull() {
  // take care of the nil!
  sz = 1;
  sum = val;
  for(int i = 0; i < 2; i++){</pre>
    if(ch[i] == &nil) continue;
    ch[i]->f = this;
    sz += ch[i]->sz;
    sum += ch[i]->sum;
  }
}
void rotate(){
  Splay *p = f;
  int d = dir();
  if (!p->isr()) p->f->setCh(this, p->dir());
  else f = p->f;
  p->setCh(ch[!d], d);
  setCh(p, !d);
  p->pull(), pull();
void update(){
  if(f != &nil) f->update();
  push();
void splay(){
  update():
  for(Splay* fa; fa = f, !isr(); rotate())
    if(!fa->isr()) (fa->dir() == dir() ? fa : this)->
        rotate();
Splay *access(Splay* q = &nil){
  splay();
  setCh(q, 1);
  pull();
  if (f != &nil) return f->access(this);
  else return q;
void root_path(){access(), splay();}
void chroot() {root_path(), swap(ch[0], ch[1]), rev = 1,
    push(), pull();}
void split(Splay* y){chroot(), y->root_path();}
void link(Splay* y){root_path(), y->chroot(), setCh(y, 1)
void cut(Splay* y) {split(y), y->push(), y->ch[0] = y->ch
    [0] - f = &nil;
Splay *get_root(){
  root_path();
```

```
auto q = this;
    for(; q->ch[0] != &nil; q = q->ch[0]) q->push();
    return q;
  Splay *lca(Splay* y){
    access(), y->root_path();
    return y->f == &nil ? &nil : y->f;
  bool conn(Splay* y){return get_root() == y->get_root();}
} Splay::nil;
2.3
      Treap
struct node {
  int data, sz;
  node *1, *r;
  node(int k) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
};
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a | | !b) return a ? a : b;
  if (rand() \% (sz(a) + sz(b)) < sz(a))
    return a \rightarrow down(), a \rightarrow r = merge(a \rightarrow r, b), a \rightarrow up(),
            a:
  return b \rightarrow down(), b \rightarrow 1 = merge(a, b \rightarrow 1), b \rightarrow up(), b;
void split(node *o, node *&a, node *&b, int k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), <math>a->up();
  else b = o, split(o \rightarrow l, a, b \rightarrow l, k), b \rightarrow up();
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - sz(o->l) - 1);
  else b = o, split2(o \rightarrow 1, a, b \rightarrow 1, k);
  o->up();
node *kth(node *o, int k) {
  if (k \le sz(o->1)) return kth(o->1, k);
  if (k == sz(o\rightarrow 1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, int key) {
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o\rightarrow 1) + 1 + Rank(o\rightarrow r, key);
  else return Rank(o->1, key);
bool erase(node *&o, int k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    delete t;
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, int k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
void interval(node *&o, int 1, int r) {
  node *a, *b, *c;
```

```
split2(o, a, b, l - 1), split2(b, b, c, r);
  // operate
 o = merge(a, merge(b, c));
    KD Tree
2.4
namespace kdt {
 int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
 yl[maxn], yr[maxn];
  point p[maxn];
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
    function<bool(const point &, const point &)> f =
      [dep](const point &a, const point &b) {
        if (dep & 1) return a.x < b.x;</pre>
        else return a.y < b.y;</pre>
    int m = (1 + r) >> 1;
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
    yl[m] = yr[m] = p[m].y;
lc[m] = build(1, m, dep + 1);
    if (~lc[m]) {
      xl[m] = min(xl[m], xl[lc[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
    }
    return m;
  bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
        q.y < y1[o] - ds || q.y > yr[o] + ds
      return false;
    return true;
  long long dist(const point &a, const point &b) {
    return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
  void dfs(
      const point &q, long long &d, int o, int dep = 0) {
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x ||
        !(dep & 1) && q.y < p[o].y) {
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    } else {
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
 void init(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
    root = build(0, v.size());
  long long nearest(const point &q) {
    long long res = 1e18;
    dfs(q, res, root);
    return res;
} // namespace kdt
```

```
struct node {
  11 v, data, sz, sum;
  node *1, *r;
  node(ll k)
    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a->r) > V(a->1)) swap(a->r, a->1);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
     Flow & Matching
3
3.1
     Dinic
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2})), bipartite matching: O(sqrt(V)E)
  struct edge {
    ll to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
    n = n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
    g[u].pb(edge{v, cap, 0, SZ(g[v])});
    g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
    dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int cur = q.front(); q.pop();
      for (auto &e : g[cur]) {
        if (dis[e.to] == -1 && e.flow != e.cap) {
          q.push(e.to);
          dis[e.to] = dis[cur] + 1;
      }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
```

g[e.to][e.rev].flow -= df;

int u = q.front();

q.pop(), inq[u] = 0;

```
return df;
        }
     }
    dis[u] = -1;
    return 0;
  11 maxflow(int _s, int _t) {
    s = _s; t = _t;
11 flow = 0, df;
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
};
    Bounded Flow
struct BoundedFlow : Dinic {
 vector<ll> tot;
 void init(int _n) {
    Dinic::init(_n + 2);
    tot.assign(n, 0);
 void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
    g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
 bool feasible() {
    11 \text{ sum } = 0;
    int vs = n - 2, vt = n - 1;
    for(int i = 0; i < n - 2; ++i)
      if(tot[i] > 0)
        add_edge(vs, i, 0, tot[i]), sum += tot[i];
      else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
        g[vs].pop_back(), g[i].pop_back();
      else if(tot[i] < 0)</pre>
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
    add_edge(_t, _s, 0, INF);
    if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    //return x + maxflow(_s, _t); // max
};
3.3 MCMF
struct MCMF { // 0-based, O(SPFA * |f|)
  struct edge {
    11 from, to, cap, flow, cost, rev;
 }:
 int n;
 int s, t; ll mx;
 //mx: maximum amount of flow
 vector<vector<edge>> g;
 vector<ll> dis, up;
 bool BellmanFord(ll &flow, ll &cost) {
    vector<edge*> past(n);
    vector<int> inq(n);
    dis.assign(n, INF); up.assign(n, 0);
    queue<int> q;
    q.push(s), inq[s] = 1;
```

up[s] = mx - flow, past[s] = 0, dis[s] = 0;

while (!q.empty()) {

```
if (!up[u]) continue;
      for (auto &e : g[u])
        if (e.flow != e.cap &&
            dis[e.to] > dis[u] + e.cost) {
          dis[e.to] = dis[u] + e.cost, past[e.to] = &e;
          up[e.to] = min(up[u], e.cap - e.flow);
          if (!inq[e.to]) inq[e.to] = 1, q.push(e.to);
    if (dis[t] == INF) return 0;
    flow += up[t], cost += up[t] * dis[t];
    for (ll i = t; past[i]; i = past[i] \rightarrow from) {
      auto &e = *past[i];
      e.flow += up[t], g[e.to][e.rev].flow -= up[t];
    }
    return 1;
  }
  pll MinCostMaxFlow(int _s, int _t) {
    s = _s, t = _t;
    11 \text{ flow} = 0, \text{ cost} = 0;
    while (BellmanFord(flow, cost));
    return pll(flow, cost);
  void init(int _n, ll _mx) {
    n = n, mx = mx;
    g.assign(n, vector<edge>());
  void add_edge(int a, int b, ll cap, ll cost) {
    g[a].pb(edge{a, b, cap, 0, cost, SZ(g[b])});
    g[b].pb(edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
};
3.4 Min Cost Circulation
struct MinCostCirculation { // 0-based, O(VE * ElogC)
  struct edge {
    ll from, to, cap, fcap, flow, cost, rev;
  int n;
  vector<edge*> past;
  vector<vector<edge>> g;
  vector<ll> dis;
  void BellmanFord(int s) {
    vector<int> inq(n);
    dis.assign(n, INF);
    queue<int> q;
    auto relax = [&](int u, ll d, edge *e) {
      if (dis[u] > d) {
        dis[u] = d, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
    };
    relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u])
        if (e.cap > e.flow)
          relax(e.to, dis[u] + e.cost, &e);
    }
  void try_edge(edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
      ++cur.flow, --g[cur.to][cur.rev].flow;
      for (int i = cur.from; past[i]; i = past[i] -> from) {
        auto &e = *past[i];
        ++e.flow, --g[e.to][e.rev].flow;
      }
    ++cur.cap;
```

};

```
void solve(int mxlg) { // mxlg >= log(max cap)
    for (int b = mxlg; b >= 0; --b) {
      for (int i = 0; i < n; ++i)
        for (auto &e : g[i])
          e.cap *= 2, e.flow *= 2;
      for (int i = 0; i < n; ++i)</pre>
        for (auto &e : g[i])
          if (e.fcap >> b & 1)
            try_edge(e);
    }
 }
 void init(int _n) {
   n = _n;
    past.assign(n, nullptr);
    g.assign(n, vector<edge>());
 void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a == b)}
    g[b].pb(edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
};
```

3.5 Gomory Hu

```
void GomoryHu(Dinic &flow) { // 0-based
  int n = flow.n;
  vector<int> par(n);
  for (int i = 1; i < n; ++i) {
    flow.reset();
    add_edge(i, par[i], flow.maxflow(i, par[i]));
    for (int j = i + 1; j < n; ++j)
        if (par[j] == par[i] && ~flow.dis[j])
        par[j] = i;
  }
}</pre>
```

3.6 Stoer Wagner Algorithm

```
struct StoerWagner { // 0-based, O(V^3)
 vector<int> vis, del;
 vector<ll> wei;
 vector<vector<ll>> edge;
 void init(int _n) {
   n = _n;
   del.assign(n, 0);
   edge.assign(n, vector<ll>(n));
 void add_edge(int u, int v, ll w) {
   edge[u][v] += w, edge[v][u] += w;
 void search(int &s, int &t) {
   vis.assign(n, 0); wei.assign(n, 0);
   s = t = -1;
   while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
       if (!del[i] && !vis[i] && mx < wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)
       if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
   }
 11 solve() {
   11 ret = INF;
   for (int i = 0, x=0, y=0; i < n-1; ++i) {
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
      for (int j = 0; j < n; ++j)
        edge[x][j] = (edge[j][x] += edge[y][j]);
   return ret;
```

```
3.7 Bipartite Matching
```

```
//min vertex cover: take all unmatched vertices in L and
    find alternating tree,
//ans is not reached in L + reached in R
// O(VE)
int n; // 1-based, max matching
int mx[maxn], my[maxn];
bool adj[maxn][maxn], vis[maxn];
bool dfs(int u) {
  if (vis[u]) return 0;
  vis[u] = 1;
  for (int v = 1; v <= n; v++) {
    if (!adj[u][v]) continue;
    if (!my[v] || (my[v] && dfs(my[v]))) {
      mx[u] = v, my[v] = u;
      return 1;
   }
  }
  return 0;
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!\sim my[e] \mid | (dis[my[e]] == dis[u] + 1 \&\& dfs(my[e])
          ])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < nl; ++i)</pre>
      if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!\sim my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
        }
    }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    }
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
    nl = _nl, nr = _nr;
    g.assign(nl, vector<int>());
  }
};
```

3.8 Kuhn Munkres Algorithm

```
struct KM { // 0-based, maximum matching, O(V^3)
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll> hl, hr, slk;
 vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    n = _n;
    // -INF for perfect matching
    w.assign(n, vector<11>(n, 0));
    pre.assign(n, 0);
    qu.assign(n, 0);
 void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
 bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!v1[x] \&\& s1k[x] >= (d = h1[x] + hr[y] - w[x])
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!v1[x] \&\& d > s1k[x]) d = s1k[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !slk[x] && !check(x)) return;
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0); hr
         .assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res:
};
```

3.9 Max Simple Graph Matching

```
struct Matching { // 0-based, O(V^3)
 queue<int> q; int n;
 vector<int> fa, s, vis, pre, match;
 vector<vector<int>> g;
 int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
   static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
     if (vis[x] == tk) return x;
      vis[x] = tk;
     x = Find(pre[match[x]]);
   }
 void Blossom(int x, int y, int 1) {
    for (; Find(x) != 1; x = pre[y]) {
     pre[x] = y, y = match[x];
     if (s[y] == 1) q.push(y), s[y] = 0;
```

```
for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
    }
  bool Bfs(int r) {
    iota(iter(fa), 0); fill(iter(s), -1);
    q = queue < int > (); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : g[x])
        if (s[u] == -1) {
          if (pre[u] = x, s[u] = 1, match[u] == n) {
            for (int a = u, b = x, last;
                b != n; a = last, b = pre[a])
              last = match[b], match[b] = a, match[a] = b;
            return true;
          q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
          int l = LCA(u, x);
          Blossom(x, u, 1); Blossom(u, x, 1);
    return false;
  Matching(int _n) : n(_n), fa(n + 1), s(n + 1), vis(n + 1)
      , pre(n + 1, n), match(n + 1, n), g(n) {}
  void add_edge(int u, int v)
  { g[u].pb(v), g[v].pb(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

3.10 Stable Marriage

```
1: Initialize m \in M and w \in W to free
 2: while \exists free man m who has a woman w to propose to do
       w \leftarrow first woman on m's list to whom m has not yet proposed
       if \exists some pair (m', w) then
4:
           if w prefers m to m' then
5:
6:
              m' \leftarrow free
7:
              (m, w) \leftarrow engaged
8:
           end if
9:
       else
10:
           (m, w) \leftarrow engaged
11:
        end if
12: end while
```

3.11 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
- 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite
- 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T

- 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
- 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with $(\cos t, cap)=(0,d(v))$
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C + K
- Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer T
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v, v \in G$ with capacity K
- 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e))$ pdd intersect(Line a, Line b) { pdd p1, p2, p3, p4;
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u, v).
 - 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G'.
- Project selection problem
- 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
- 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \ge 0$$

4 Geometry

4.1 Geometry Template

```
using ld = 11;
using pdd = pair<ld, ld>;
using Line = pair<pdd, pdd>;
#define X first
#define Y second
// ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){    return v > eps ? 1 : ( v < -eps ? -1 : 0)
    ; }
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
```

```
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)) <=
bool seg_intersect(Line a, Line b){
  pdd p1, p2, p3, p4;
  tie(p1, p2) = a; tie(p3, p4) = b;
  if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1, p2)
      || btw(p4, p1, p2))
    return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&
    ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
  pdd p1, p2, p3, p4;
  tie(p1, p2) = a; tie(p3, p4) = b;
  ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 -
     p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) /
    abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1,
    pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq), dot(dp
  return q0 + pdd(cross(r - p0, num), dot(r - p0, num)) /
      abs2(dp);
} // from line p0--p1 to q0--q1, apply to r
```

4.2 Convex Hull

```
vector<int> getConvexHull(vector<pdd>& pts){
 vector<int> id(SZ(pts));
  iota(iter(id), 0);
  sort(iter(id), [&](int x, int y){ return pts[x] < pts[y];</pre>
       });
  vector<int> hull;
  for(int tt = 0; tt < 2; tt++){
    int sz = SZ(hull);
    for(int j : id){
      pdd p = pts[j];
      while (SZ(hull) - sz >= 2 \&\&
          cross(pts[hull.back()] - pts[hull[SZ(hull) - 2]],
            p - pts[hull[SZ(hull) - 2]]) <= 0)
        hull.pop_back();
      hull.pb(j);
    hull.pop_back();
    reverse(iter(id));
  return hull;
```

4.3 Minimum Enclosing Circle

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
   random_shuffle(iter(pts));
   pdd c = pts[0];
   ld r = 0;
   for(int i = 1; i < SZ(pts); i++){
      if(abs(pts[i] - c) <= r) continue;
      c = pts[i]; r = 0;
   for(int j = 0; j < i; j++){
      if(abs(pts[j] - c) <= r) continue;
      c = (pts[i] + pts[j]) / 2;</pre>
```

if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)

return ori(a.X, a.Y, b.Y) < 0;</pre>

return cmp(a.Y - a.X, b.Y - b.X, 0);

```
r = abs(pts[i] - c);
      for(int k = 0; k < j; k++){
                                                                  deque<Line> dq(1, arr[0]);
        if(abs(pts[k] - c) > r)
                                                                  for (auto p : arr) {
          tie(c, r) = circumcenter(pts[i], pts[j], pts[k]);
                                                                    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) == -1)
                                                                      continue:
    }
                                                                    while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.back
                                                                        ()))
  return {c, r};
                                                                      dq.pop_back();
                                                                    while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
                                                                      dq.pop_front();
                                                                    dq.pb(p);
    Minkowski Sum
                                                                  while (SZ(dq) >= 3 \&\& !isin(dq[0], dq[SZ(dq) - 2], dq.
                                                                      back()))
void reorder_poly(vector<pdd>& pnts){
                                                                    dq.pop_back();
                                                                  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
                                                                    dq.pop_front();
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn].Y</pre>
                                                                  return vector<Line>(iter(dq));
        && pnts[i].X < pnts[mn].X))
  rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
                                                                4.7 Dynamic Convex Hull
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
  reorder poly(P);
                                                                struct Line{
  reorder_poly(Q);
                                                                  11 a, b, 1 = MIN, r = MAX;
  int psz = P.size();
                                                                  Line(ll a, ll b): a(a), b(b) {}
  int qsz = Q.size();
                                                                  11 operator()(11 x) const{
  P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
                                                                    return a * x + b;
  vector<pdd> ans;
  int i = 0, j = 0;
                                                                  bool operator<(Line b) const{</pre>
  while(i < psz || j < qsz){
                                                                    return a < b.a;
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1] - P[i], Q[j + 1] - Q[j]));
                                                                  bool operator<(11 b) const{</pre>
    if(t >= 0) i++;
                                                                    return r < b;
    if(t <= 0) j++;
                                                                };
  return ans;
                                                                11 iceil(l1 a, l1 b){
}
                                                                  if(b < 0) a *= -1, b *= -1;
                                                                  if(a > 0) return (a + b - 1) / b;
                                                                  else return a / b;
4.5 Polar Angle Comparator
// -1: a // b (if same), 0/1: a < b
                                                                11 intersect(Line a, Line b){
int cmp(pll a, pll b, bool same = true){
                                                                  return iceil(a.b - b.b, b.a - a.a);
#define is_neg(k) (sgn(k.Y) < 0 \mid \mid (sgn(k.Y) == 0 && sgn(k.
    X) < 0)
  int A = is_neg(a), B = is_neg(b);
                                                                struct DynamicConvexHull{
  if(A != B)
                                                                  multiset<Line, less<>> ch;
    return A < B;
  if(sgn(cross(a, b)) == 0)
                                                                  void add(Line ln){
    return same ? abs2(a) < abs2(b) : -1;</pre>
                                                                    auto it = ch.lower_bound(ln);
  return sgn(cross(a, b)) > 0;
                                                                    while(it != ch.end()){
                                                                      Line tl = *it;
                                                                      if(tl(tl.r) <= ln(tl.r)){
                                                                        it = ch.erase(it);
4.6 Half Plane Intersection
                                                                      else break;
// from 8BQube
                                                                    }
pll area_pair(Line a, Line b)
                                                                    auto it2 = ch.lower_bound(ln);
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a.X,
                                                                    while(it2 != ch.begin()){
    b.Y - a.X)); }
                                                                      Line tl = *prev(it2);
bool isin(Line 10, Line 11, Line 12) {
                                                                      if(tl(tl.1) <= ln(tl.1)){</pre>
  // Check inter(l1, l2) strictly in l0
                                                                        it2 = ch.erase(prev(it2));
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
                                                                      else break;
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128) a02Y * a12X - (__int128) a02X * a12Y >
                                                                    it = ch.lower_bound(ln);
      0; // C^4
                                                                    if(it != ch.end()){
                                                                      Line tl = *it;
/* Having solution, check size > 2 */
                                                                      if(tl(tl.1) >= ln(tl.1)) ln.r = tl.1 - 1;
/* --^-- Line.X --^-- Line.Y --^-- */
                                                                      else{
vector<Line> halfPlaneInter(vector<Line> arr) {
                                                                        11 pos = intersect(ln, tl);
  sort(iter(arr), [&](Line a, Line b) -> int {
                                                                        tl.1 = pos;
```

ln.r = pos - 1;

ch.erase(it);

ch.insert(t1);

}

```
}
    it2 = ch.lower bound(ln);
    if(it2 != ch.begin()){
      Line tl = *prev(it2);
      if(tl(tl.r) >= ln(tl.r)) ln.l = tl.r + 1;
      else{
        11 pos = intersect(t1, ln);
        tl.r = pos - 1;
        ln.1 = pos;
        ch.erase(prev(it2));
        ch.insert(t1);
    if(ln.l <= ln.r) ch.insert(ln);</pre>
 11 query(11 pos){
    auto it = ch.lower_bound(pos);
    if(it == ch.end()) return 0;
    return (*it)(pos);
};
4.8
      3D Point
  double x, y, z;
      , y(_y), z(_z){}
```

```
// Copy from 8BQube
struct Point {
 Point(double _x = 0, double _y = 0, double _z = 0): x(_x)
 Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator - (Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z); }
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z); }
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x - p1.
    x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); } double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-pi,
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0, pi]
double theta(Point p) { return atan2(sqrt(p.x * p.x + p.y *
     p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
\{ return (a + b + c + d) / 4; \}
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
 Point e1 = b - a;
 Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
 e2 = e2 / abs(e2);
 Point p = u - a;
 return pdd(dot(p, e1), dot(p, e2));
Point rotate_around(Point p, double angle, Point axis) {
  double s = sin(angle), c = cos(angle);
 Point u = axis / abs(axis);
  return u * dot(u, p) * (1 - c) + p * c + cross(u, p) * s;
```

4.9 ConvexHull3D

```
struct convex_hull_3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector<Face> res;
vector<Point> P;
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn(
      abs2(P[0] - p)) != 0; }));
  swap(P[2], *find_if(iter(P), [&](auto p) { return sgn(
      abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3],\ *find\_if(iter(P),\ [\&](auto\ p)\ \{\ return\ sgn(
      volume(P[0], P[1], P[2], p)) != 0; }));
  vector<vector<int>> flag(n, vector<int>(n));
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
    vector<Face> next;
    for (auto f : res) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]));
      if (d <= 0) next.pb(f);
      int ff = (d > 0) - (d < 0);
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a] = ff
    for (auto f : res) {
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)
          next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    }
    res = next;
  }
bool same(Face s, Face t) {
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
      return 0;
   \textbf{if } (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) \; != \; 0) \\
      return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
      return 0;
  return 1;
int polygon_face_num() {
  int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin() + i, [&](Face g
        ) { return same(res[i], g); });
  return ans;
double get_volume() {
  double ans = 0;
  for (auto f : res)
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c]);
  return fabs(ans / 6);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1.z)
      * (p3.y - p1.y);
                    - p1.z) * (p3.x - p1.x) - (p2.x - p1.x)
  double b = (p2.z)
      * (p3.z - p1.z);
  double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1.y)
      * (p3.x - p1.x);
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a * a
```

+ b * b + c * c);

} tool:

```
}
};
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point
// (0, 0, inf) to avoid degenerate case
```

4.10 Circle Operations

```
// from 8BQube
const double PI=acos(-1);
vector<pdd> circleLineIntersection(pdd c, double r, pdd a,
    pdd b) {
  pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a);
  double s = cross(b - a, c - a), h2 = r * r - s * s / abs2
      (b - a);
  if (sgn(h2) < 0) return {};</pre>
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return \{p - h, p + h\};
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r)
        -h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double areaPolyCircle(const vector<pdd> poly,const pdd &0,
    const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)\%SZ(poly)]-0,r)*ori(0,poly)
        [i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
  double r1 = a.R, r2 = b.R, d2 = abs2(o1 - o2), d = sqrt(
      d2);
  if(d < max(r1, r2) - min(r1, r2) | | d > r1 + r2) return
      0;
  pdd u = (o1 + o2) * 0.5 + (o1 - o2) * ((r2 * r2 - r1 * r1))
      ) / (2 * d2));
  double A = sqrt((r1 + r2 + d) * (r1 - r2 + d) * (r1 + r2)
      - d) * (-r1 + r2 + d));
  pdd v = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A / (2 * d2);
  p1 = u + v, p2 = u - v;
  return 1;
vector<Line> CCtang( const Cir& c1 , const Cir& c2 , int
    sign1 ){
  vector<Line> ret;
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
 for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
      v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
```

```
pdd p2 = c2.0 + n * (c2.R * sign1);
if (sgn(p1.X - p2.X) == 0 and
         sgn(p1.Y - p2.Y) == 0)
    p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
}
return ret;
}
```

```
4.11 Delaunay Triangulation
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle. */
struct Edge {
  int id; // oidx[id]
  list<Edge>::iterator twin;
  Edge(int _id = 0):id(_id) {}
};
struct Delaunay { // 0-base
  int n, oidx[N];
  list<Edge> head[N]; // result udir. graph
  pll p[N];
  void init(int _n, pll _p[]) {
    n = _n, iota(oidx, oidx + n, 0);
    for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(oidx, oidx + n, [&](int a, int b)
    { return _p[a] < _p[b]; });
    for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];
    divide(0, n - 1);
  void addEdge(int u, int v) {
    head[u].push_front(Edge(v));
    head[v].push_front(Edge(u));
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
  void divide(int 1, int r) {
    if (l == r) return;
    if (1 + 1 == r) return addEdge(1, 1 + 1);
    int mid = (1 + r) >> 1, nw[2] = \{1, r\};
    divide(l, mid), divide(mid + 1, r);
    auto gao = [&](int t) {
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      for (auto it : head[nw[t]]) {
        int v = ori(pt[1], pt[0], p[it.id]);
        if (v > 0 \mid | (v == 0 \&\& abs2(pt[t ^ 1] - p[it.id])
            < abs2(pt[1] - pt[0])))</pre>
          return nw[t] = it.id, true;
      return false;
    while (gao(0) || gao(1));
    addEdge(nw[0], nw[1]); // add tangent
    while (true) {
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      int ch = -1, sd = 0;
      for (int t = 0; t < 2; ++t)
          for (auto it : head[nw[t]])
              if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
                    -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
                   id])))
      ch = it.id, sd = t;
if (ch == -1) break; // upper common tangent
      for (auto it = head[nw[sd]].begin(); it != head[nw[sd
          ]].end(); )
        if (seg_strict_intersect(pt[sd], p[it->id], pt[sd ^
             1], p[ch]))
          head[it->id].erase(it->twin), head[nw[sd]].erase(
              it++);
        else ++it;
      nw[sd] = ch, addEdge(nw[0], nw[1]);
    }
  }
```

4.12 Voronoi Diagram

4.13 Polygon Union

```
// from 8BQube
ld rat(pll a, pll b) {
 return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
 // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
  1d res = 0;
 for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {
      pll A = p[a], B = p[(a + 1) % SZ(p)];
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
          pll C = q[b], D = q[(b + 1) \% SZ(q)];
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            ld sa = cross(D - C, A - C), sb = cross(D - C,
                B - C);
            segs.pb(sa / (sa - sb), sgn(sc - sd));
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D - C</pre>
              )) > 0) {
            segs.pb(rat(C - A, B - A), 1);
            segs.pb(rat(D - A, B - A), -1);
          }
        }
      sort(iter(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
      res += cross(A, B) * sum;
 return res / 2;
```

4.14 Tangent Point to Convex Hull

```
// from 8BQube
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
      { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

5 Graph

5.1 Block Cut Tree

```
struct BCC{
  vector<int> v, e, cut;
struct BlockCutTree{ // 0-based, allow multi edges but not
    allow loops
  int n, m, cnt = 0;
  // n:|V|, m:|E|, cnt:|bcc|
  vector<BCC> bcc;
  vector<vector<pii>>> g; // original graph
  vector<pii> edges; // 0-based
  vector<vector<int>> vbcc;
  // vbcc[i] = BCCs containing vertex i, vbcc[i].size()>1
      iff i is an articulation
  vector<int> ebcc;
 // edge i is a bridge iff bcc[ebcc[i]].e.size() == 1
 // block cut tree:
 // adj[bcc i]: bcc[i].cut
 // adj[cut i]: vbcc[i]
  BlockCutTree(int _n, vector<pii> _edges):
      n(_n), m(SZ(_edges)), g(_n), edges(_edges), vbcc(_n),
           ebcc(SZ(_edges)){
    for(int i = 0; i < m; i++){</pre>
      auto [u, v] = edges[i];
      g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
  void build(){
    vector<int> in(n, -1), low(n, -1);
    vector<vector<int>> up(n);
    vector<int> stk;
    int ts = 0;
    auto _dfs = [&](auto dfs, int now, int par, int pe) ->
      if(pe != -1) up[now].pb(pe);
      in[now] = low[now] = ts++;
      stk.pb(now);
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
        if(in[v] != -1){
          if(in[v] < in[now]) up[now].pb(e);</pre>
          low[now] = min(low[now], in[v]);
          continue;
        dfs(dfs, v, now, e);
        low[now] = min(low[now], low[v]);
      if((now != par && low[now] >= in[par]) || (now == par
           && SZ(g[now]) == 0)){
        bcc.pb();
        while(true){
          int v = stk.back();
          stk.pop_back();
          vbcc[v].pb(cnt);
          bcc[cnt].v.pb(v);
          for(int e : up[v]){
            ebcc[e] = cnt;
            bcc[cnt].e.pb(e);
          if(v == now) break;
        if(now != par){
          vbcc[par].pb(cnt);
          bcc[cnt].v.pb(par);
        cnt++;
      }
    };
    for(int i = 0; i < n; i++){</pre>
      if(in[i] == -1) _dfs(_dfs, i, i, -1);
    for(int i = 0; i < cnt; i++)</pre>
```

```
for(int j : bcc[i].v)
        if(SZ(vbcc[j]) > 1) bcc[i].cut.pb(j);
 }
      2-SAT
5.2
struct SAT{ // 0-based, [n, 2n) is neg of [0, n)
 int n;
  vector<vector<int>> g, rg;
 bool ok = true;
 vector<bool> ans;
 void init(int _n){
    n = _n;
    g.resize(2 * n);
    rg.resize(2 * n);
    ans.resize(n);
  int neg(int v){
    return v < n ? v + n : v - n;
 void addEdge(int u, int v){
    g[u].pb(v);
    rg[v].pb(u);
 void addClause(int a, int b){
    addEdge(neg(a), b);
    addEdge(neg(b), a);
 void build(){
    vector<bool> vst(2 * n, false);
    vector<int> tmp, scc(2 * n, -1);
    int cnt = 1;
    function < void(int) > dfs = [&](int now){
      vst[now] = true;
      for(int i : rg[now]){
        if(vst[i]) continue;
        dfs(i);
      tmp.pb(now);
    };
    for(int i = 0; i < 2 * n; i++){
      if(!vst[i]) dfs(i);
    reverse(all(tmp));
    function<void(int, int)> dfs2 = [&](int now, int id){
      scc[now] = id;
      for(int i : g[now]){
        if(scc[i] != -1) continue;
        dfs2(i, id);
      }
    };
    for(int i : tmp){
      if(scc[i] == -1) dfs2(i, cnt++);
    debug(scc);
    for(int i = 0; i < n; i++){</pre>
      if(scc[i] == scc[neg(i)]){
        ok = false;
        return:
      if(scc[i] < scc[neg(i)]) ans[i] = true;</pre>
      else ans[i] = false;
 }
};
     Dominator Tree
// copy from 8BQube
struct dominator_tree { // 1-base
```

```
// copy from 8BQube
struct dominator_tree { // 1-base
  vector<int> G[N], rG[N];
  int n, pa[N], dfn[N], id[N], Time;
```

```
int semi[N], idom[N], best[N];
  vector<int> tree[N]; // dominator_tree
  void init(int n) {
    n = _n;
for (int i = 1; i <= n; ++i)</pre>
      G[i].clear(), rG[i].clear();
  void add_edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
      if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  int find(int y, int x) {
    if (y <= x) return y;</pre>
    int tmp = find(pa[y], x);
    if (semi[best[y]] > semi[best[pa[y]]])
      best[y] = best[pa[y]];
    return pa[y] = tmp;
  void tarjan(int root) {
    Time = 0:
    for (int i = 1; i <= n; ++i) {
      dfn[i] = idom[i] = 0;
      tree[i].clear();
      best[i] = semi[i] = i;
    dfs(root);
    for (int i = Time; i > 1; --i) {
      int u = id[i];
      for (auto v : rG[u])
        if (v = dfn[v]) {
          find(v, i);
          semi[i] = min(semi[i], semi[best[v]]);
      tree[semi[i]].pb(i);
      for (auto v : tree[pa[i]]) {
        find(v, pa[i]);
        idom[v] =
          semi[best[v]] == pa[i] ? pa[i] : best[v];
      tree[pa[i]].clear();
    for (int i = 2; i <= Time; ++i) {</pre>
      if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
      tree[id[idom[i]]].pb(id[i]);
    }
  }
};
      Virtual Tree
5.4
// copy from 8BQube
vector<int> vG[N];
int top, st[N];
int vrt = -1;
void insert(int u) {
  if (top == -1) return st[++top] = vrt = u, void();
  int p = LCA(st[top], u);
    if(dep[vrt] > dep[p]) vrt = p;
  if (p == st[top]) return st[++top] = u, void();
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
    vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
```

void reset(int u) {

vG[u].clear();

for (int i : vG[u]) reset(i);

```
void solve(vector<int> &v) {
 top = -1;
 sort(ALL(v),
      [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
 for (int i : v) insert(i);
 while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
 // do something
 reset(vrt);
      Directed Minimum Spanning Tree
const 11 INF = LLONG_MAX;
struct edge{
 int u = -1, v = -1;
 11 w = INF;
 int id = -1;
// 0-based, E[i].id = i
bool DMST(int n, vector<edge> &E, int root, vector<edge> &
    sol){
 vector<int> id(n), vis(n);
 vector<edge> in(n);
 for(edge e : E)
   if(e.u != e.v && e.w < in[e.v].w && e.v != root)</pre>
      in[e.v] = e;
 for(int i = 0; i < n; i++)</pre>
   if(i != root && in[i].u == -1) return false; // no sol
  int cnt = 0;
 fill(iter(id), -1); fill(iter(vis), -1);
 for(int u = 0; u < n; u++){}
   int v = u;
   while(vis[v] != u && id[v] == -1 && in[v].u != -1)
      vis[v] = u, v = in[v].u;
   if(v != root && id[v] == -1){
      for(int x = in[v].u; x != v; x = in[x].u)
       id[x] = cnt;
      id[v] = cnt++;
   }
 if(!cnt) return sol = in, true; // no cycle
 for(int u = 0; u < n; u++)</pre>
   if(id[u] == -1) id[u] = cnt++;
 vector<edge> nE;
 for(int i = 0; i < SZ(E); i++){</pre>
   edge tmp = E[i];
    tmp.u = id[tmp.u], tmp.v = id[tmp.v];
   if(in[E[i].v].w != INF) tmp.w -= in[E[i].v].w;
   nE.pb(tmp);
 vector<edge> tsol;
 if(!DMST(cnt, nE, id[root], tsol)) return false;
  sol.resize(n);
 for(int i = 0; i < cnt; i++){</pre>
   if(i == id[root]) continue;
   int t = tsol[i].id;
   sol[E[t].v] = E[t];
 for(int i = 0; i < n; i++)</pre>
   if(sol[i].id == -1) sol[i] = in[i];
 return true;
5.6 Fast DMST
struct Edge { int a, b; ll w; };
struct Node { /// Lazy skew heap node
 Edge key;
 Node *1, *r;
 ll delta;
 void prop() {
   key.w += delta;
   if (1) 1->delta += delta;
```

if (r) r->delta += delta;

```
delta = 0;
  Edge top() { prop(); return key; }
};
Node *merge(Node *a, Node *b) {
  if (!a || !b) return a ?: b;
  a->prop(), b->prop();
  if (a->key.w > b->key.w) swap(a, b);
  swap(a->1, (a->r = merge(b, a->r)));
  return a:
void pop(Node*\& a) \{ a->prop(); a = merge(a->l, a->r); \}
pair<11, vi> dmst(int n, int r, vector<Edge>& g) {
  RollbackUF uf(n); // need to implement this
  vector<Node*> heap(n);
  for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e
      });
  11 \text{ res} = 0;
  vi seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<Edge> Q(n), in(n, \{-1,-1\}), comp;
  deque<tuple<int, int, vector<Edge>>> cycs;
  rep(s,0,n) {
    int u = s, qi = 0, w;
    while (seen[u] < 0) {
      if (!heap[u]) return {-1,{}};
      Edge e = heap[u]->top();
      heap[u]->delta -= e.w, pop(heap[u]);
      Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if (seen[u] == s) { /// found cycle, contract
        Node* cyc = 0;
        int end = qi, time = uf.time();
        do cyc = merge(cyc, heap[w = path[--qi]]);
        while (uf.join(u, w));
        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
        cycs.push_front({u, time, {&Q[qi], &Q[end]}});
      }
    }
    rep(i,0,qi) in[uf.find(Q[i].b)] = Q[i];
  for (auto& [u,t,comp] : cycs) { // restore sol (optional)
    uf.rollback(t);
    Edge inEdge = in[u];
    for (auto& e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
  }
  rep(i,0,n) par[i] = in[i].a;
  return {res, par};
5.7
     Vizing
// find D+1 edge coloring of a graph with max deg D
struct vizing { // returns edge coloring in adjacent matrix
     G. 1 - based
  const int N = 105;
  int C[N][N], G[N][N], X[N], vst[N], n; // ans: G[i][j]
  void init(int _n) { n = _n; // n = |V|+1
    for (int i = 0; i <= n; ++i)
      for (int j = 0; j <= n; ++j)
        C[i][j] = G[i][j] = 0;
  void solve(vector<pii> &E) {
    auto update = [&](int u)
    { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
    auto color = [&](int u, int v, int c) {
      int p = G[u][v];
      G[u][v] = G[v][u] = c;
      C[u][c] = v, C[v][c] = u;
      C[u][p] = C[v][p] = 0;
      if (p) X[u] = X[v] = p;
      else update(u), update(v);
      return p:
```

```
};
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
    if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill_n(X + 1, n, 1);
  for (int t = 0; t < SZ(E); ++t) {
    int u = E[t].X, v0 = E[t].Y, v = v0, c0 = X[u], c =
        c0, d;
    vector<pii> L;
    fill_n(vst + 1, n, 0);
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0; --a)
          c = color(u, L[a].X, c);
      else if (!C[u][d]) for (int a = SZ(L) - 1; a >= 0;
          --a) color(u, L[a].X, L[a].Y);
      else if (vst[d]) break;
      else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
      if (int a; C[u][c0]) {
        for (a = SZ(L) - 2; a >= 0 && L[a].Y != c; --a);
        for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
      }
      else --t;
}
```

Maximum Clique

```
struct MaxClique { // fast when N <= 100</pre>
 bitset<N> G[N], cs[N];
 int ans, sol[N], q, cur[N], d[N], n;
 void init(int _n) {
   n = _n;
   for (int i = 0; i < n; ++i) G[i].reset();</pre>
 void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
 void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
   if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(ALL(r), [\&](int x, int y) \{ return d[x] > d[y];
   vector<int> c(SZ(r));
   int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k].
           _Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
   dfs(r, c, l + 1, mask);
 void dfs(vector<int> &r, vector<int> &c, int 1, bitset<N>
       mask) {
   while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
```

```
cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
};
```

5.9 Number of Maximal Clique

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
        if (g[v][some[d][j]])
          some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)
        if (g[v][none[d][j]])
          none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
    }
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
};
```

5.10 Minimum Mean Cycle

```
// from 8BQube
11 road[N][N]; // input here
struct MinimumMeanCycle {
 11 dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
    for (int i = 2; i <= L; ++i)
      for (int k = 0; k < n; ++k)
        for (int j = 0; j < n; ++j)
          dp[i][j] =
            min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
      if (dp[L][i] >= INF) continue;
      11 ta = 0, tb = 1;
      for (int j = 1; j < n; ++j)
        if (dp[j][i] < INF \&\&
          ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
```

```
ta = dp[L][i] - dp[j][i], tb = L - j;
if (ta == 0) continue;
if (a == -1 || a * tb > ta * b) a = ta, b = tb;
}
if (a != -1) {
    ll g = __gcd(a, b);
    return pll(a / g, b / g);
}
return pll(-1LL, -1LL);
}
void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
}
};</pre>
```

5.11 Minimum Steiner Tree

```
// from 8BQube
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
  int vcost[N]; // the cost of vertexs
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
  void add_edge(int ui, int vi, int wi) {
    dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j)
          dst[i][j] =
            min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
             submsk = (submsk - 1) \& msk)
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
              vcost[i]);
      for (int i = 0; i < n; ++i) {
        tdst[i] = INF;
        for (int j = 0; j < n; ++j)
          tdst[i] =
            min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    int ans = INF;
    for (int i = 0; i < n; ++i)</pre>
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
};
```

6 Math

return g0;

}

```
Extended Euclidean Algorithm
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
void extgcd(ll x,ll y,ll &g,ll &a,ll &b) {
  if (y == 0) g=x,a=1,b=0;
  else extgcd(y,x%y,g,b,a),b-=(x/y)*a;
6.2 Floor & Ceil
ll ifloor(ll a,ll b){
  return a / b - (a % b && (a < 0) ^ (b < 0));
11 iceil(ll a,ll b){
  return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre
// the Jacobi symbol is a generalization of the Legendre
// such that the bottom doesn't need to be prime.
// (n|p) -> same as Legendre
// (n|ab) = (n|a)(n|b)
// work with long long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
   a %= m;
   if (a == 0) return 0;
    const int r = __builtin_ctz(a);
   if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
   a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  return s;
// 0: a == 0
// -1: a isn't a quad res of p
// else: return X with X^2 % p == a
// doesn't work with long long
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0) return 0;
  if (jc == -1) return -1;
  int b, d;
  for (;;) {
   b = rand() % p;
    d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p))
           % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) %
    f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  }
```

6.4 Simplex

```
#pragma once
typedef double T; // long double, Rational, double + mod<P
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1/.0;
#define MP make_pair
#define ltj(X) if(s == -1 || MP(X[j],N[j]) < MP(X[s],N[s]))
struct LPSolver {
 int m, n;
 vi N, B;
 vvd D;
 LPSolver(const vvd& A, const vd& b, const vd& c) :
    m(sz(b)), n(sz(c)), N(n+1), B(m), D(m+2), vd(n+2)) {
      rep(i,0,m) \ rep(j,0,n) \ D[i][j] = A[i][j];
      rep(i,0,m) { B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[
          i];}
      rep(j,0,n) \{ N[j] = j; D[m][j] = -c[j]; \}
      N[n] = -1; D[m+1][n] = 1;
 void pivot(int r, int s) {
    T *a = D[r].data(), inv = 1 / a[s];
    rep(i,0,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
      T *b = D[i].data(), inv2 = b[s] * inv;
      rep(j,0,n+2) b[j] -= a[j] * inv2;
      b[s] = a[s] * inv2;
    rep(j,0,n+2) if (j != s) D[r][j] *= inv;
    rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
 bool simplex(int phase) {
    int x = m + phase - 1;
    for (;;) {
      int s = -1:
      rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
      if (D[x][s] >= -eps) return true;
      int r = -1;
      rep(i,0,m) {
        if (D[i][s] <= eps) continue;</pre>
        if (r == -1 || MP(D[i][n+1] / D[i][s], B[i])
            < MP(D[r][n+1] / D[r][s], B[r])) r = i;
      if (r == -1) return false;
      pivot(r, s);
 T solve(vd &x) {
    int r = 0;
    rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] < -eps) {
      pivot(r, n);
      if (!simplex(2) || D[m+1][n+1] < -eps) return -inf;</pre>
      rep(i,0,m) if (B[i] == -1) {
        int s = 0;
        rep(j,1,n+1) ltj(D[i]);
        pivot(i, s);
     }
    bool ok = simplex(1); x = vd(n);
    rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
    return ok ? D[m][n+1] : inf;
};
```

6.5 Floor Sum

```
// from 8BQube
11 floor_sum(ll n, ll m, ll a, ll b) {
  assert(m);
  if(m < 0) return -floor_sum(n, -m, a, b-m-1);</pre>
  11 \text{ ans} = 0;
  if (a >= m)
   ans += (n - 1) * n * (a / m) / 2, a %= m;
  if (b >= m)
    ans += n * (b / m), b %= m;
  11 y_max = (a * n + b) / m, x_max = (y_max * m - b);
 if (y_max == 0) return ans;
 ans += (n - (x_max + a - 1) / a) * y_max;
  ans += floor_sum(y_max, a, m, (a - x_max % a) % a);
 return ans;
// sum^{n-1}_0 floor((a * i + b) / m) in log(n + m + a + b)
    )
```

6.6 DiscreteLog

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered_map<int, int> p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
    y = 1LL * y * x % m;
    b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {</pre>
    s = 1LL * s * b % m;
    if (p.find(s) != p.end()) return i + kStep - p[s];
  }
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
   if (s == y) return i;
s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p; //returns: x^p = y \pmod{m}
}
```

6.7 Miller Rabin & Pollard Rho

```
// n < 4,759,123,141
                         3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
11 mul(11 a, 11 b, 11 n){
 return (__int128)a * b % n;
bool Miller_Rabin(ll a, ll n) {
 if ((a = a % n) == 0) return 1;
  if (n % 2 == 0) return n == 2;
 11 \text{ tmp} = (n - 1) / ((n - 1) & (1 - n));
  for (; tmp; tmp >>= 1, a = mul(a, a, n))
   if (tmp & 1) x = mul(x, a, n);
  if (x == 1 || x == n - 1) return 1;
  while (--t)
   if ((x = mul(x, x, n)) == n - 1) return 1;
  return 0;
bool prime(ll n){
  vector<11> tmp = {2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
```

```
for(ll i : tmp)
    if(!Miller_Rabin(i, n)) return false;
  return true;
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2], void
      ();
  11 x = 2, y = 2, d = 1, p = 1;
#define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return:
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
    d = gcd(abs(x - y), n);
      XOR Basis
const int digit = 60; // [0, 2^digit)
struct Basis{
  int total = 0, rank = 0;
  vector<ll> b;
  Basis(): b(digit) {}
  bool add(ll v){ // Gauss Jordan Elimination
    total++;
    for(int i = digit - 1; i >= 0; i--){
      if(!(1LL << i & v)) continue;</pre>
      if(b[i] != 0){
        v ^= b[i];
        continue;
      for(int j = 0; j < i; j++)</pre>
        if(1LL << j & v) v ^= b[j];</pre>
      for(int j = i + 1; j < digit; j++)</pre>
        if(1LL << i & b[j]) b[j] ^= v;
      b[i] = v;
      rank++;
      return true;
    return false;
  11 \text{ getmax}(11 \text{ x = 0})
    for(ll i : b) x = max(x, x ^ i);
    return x;
  11 \text{ getmin}(11 \text{ x} = 0)
    for(ll i : b) x = min(x, x ^ i);
    return x;
  bool can(ll x){
    return getmin(x) == 0;
  11 kth(11 k){ // kth smallest, 0-indexed
    vector<ll> tmp;
    for(ll i : b) if(i) tmp.pb(i);
    11 \text{ ans} = 0;
    for(int i = 0; i < SZ(tmp); i++)</pre>
      if(1LL << i & k) ans ^= tmp[i];</pre>
    return ans;
  }
};
```

6.9 Linear Equation

```
vector<int> RREF(vector<vector<11>> &mat){
  int N = mat.size(), M = mat[0].size();
  int rk = 0;
```

```
vector<int> cols;
  for (int i = 0;i < M;i++) {</pre>
    int cnt = -1;
    for (int j = N-1; j >= rk; j--)
      if(mat[j][i] != 0) cnt = j;
    if(cnt == -1) continue;
    swap(mat[rk], mat[cnt]);
    11 lead = mat[rk][i];
    for (int j = 0;j < M;j++) mat[rk][j] /= lead;</pre>
    for (int j = 0; j < N; j++) {
      if(j == rk) continue;
      11 tmp = mat[j][i];
      for (int k = 0; k < M; k++)
        mat[j][k] -= mat[rk][k] * tmp;
    cols.pb(i);
    rk++;
  return cols;
struct LinearEquation{
  bool ok:
  vector<11> par; //particular solution (Ax = b)
  vector<vector<ll>> homo; //homogenous (Ax = 0)
  vector<vector<ll>> rref;
  //first M columns are matrix A
  //last column of eq is vector b
  void solve(const vector<vector<ll>>> &eq){
    int M = (int)eq[0].size() - 1;
    rref = eq;
    auto piv = RREF(rref);
    int rk = piv.size();
    if(piv.size() && piv.back() == M){
      ok = 0; return;
    ok = 1;
    par.resize(M);
    vector<bool> ispiv(M);
    for (int i = 0;i < rk;i++) {</pre>
      par[piv[i]] = rref[i][M];
      ispiv[piv[i]] = 1;
    for (int i = 0; i < M; i++) {
      if (ispiv[i]) continue;
      vector<ll> h(M);
      h[i] = 1;
      for (int j = 0;j < rk;j++) h[piv[j]] = -rref[j][i];</pre>
      homo.pb(h);
    }
 }
}
6.10 Chinese Remainder Theorem
pll solve_crt(ll x1, ll m1, ll x2, ll m2){
  ll g = gcd(m1, m2);
  if ((x2 - x1) % g) return {0, 0}; // no sol
  m1 /= g; m2 /= g;
```

6.11 Sqrt Decomposition

```
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){
   int x = ifloor(n, l);
   r = ifloor(n, x);
}
// for all i in [l, r], ceil(n / i) = x</pre>
```

```
for(int l, r = n; r >= 1; r = l - 1){
 int x = iceil(n, r);
 l = iceil(n, x);
```

Misc

7.1 Cyclic Ternary Search

```
/* bool pred(int a, int b);
f(0) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
 int l = 0, r = n; bool rv = pred(1, 0);
 while (r - l > 1) {
    int m = (1 + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else l = m;
  return pred(1, r % n) ? 1 : r % n;
```

Matroid

我們稱一個二元組 $M = (E, \mathcal{I})$ 為一個擬陣,其中 $\mathcal{I} \subseteq 2^E$ 為 E 的子集所形成的 非空集合,若:

- 若 $S \in \mathcal{I}$ 以及 $S' \subsetneq S$,則 $S' \in \mathcal{I}$
- 對於 $S_1, S_2 \in \mathcal{I}$ 滿足 $|S_1| < |S_2|$,存在 $e \in S_2 \setminus S_1$ 使得 $S_1 \cup \{e\} \in \mathcal{I}$ 除此之外,我們有以下的定義:
- 位於 \mathcal{I} 中的集合我們稱之為獨立集(independent set),反之不在 \mathcal{I} 中的我們 稱為相依集(dependent set)
- 極大的獨立集為基底(base)、極小的相依集為廻路(circuit)
- 一個集合 Y 的秩 (rank) r(Y) 為該集合中最大的獨立子集,也就是 r(Y) = $\max\{|X| \mid X \subseteq Y \perp \exists X \in \mathcal{I}\}$

性質:

- 1. $X \subseteq Y \land Y \in \mathcal{I} \implies X \in \mathcal{I}$
- 2. $X \subseteq Y \land X \notin \mathcal{I} \implies Y \notin \mathcal{I}$
- 3. 若 B 與 B' 皆是基底且 $B \subseteq B'$,則 B = B'
- 若 C 與 C' 皆是迴路且 $C \subseteq C'$,則 C = C'
- 4. $e \in E \land X \subseteq E \implies r(X) \le r(X \cup \{e\}) \le r(X) + 1$ i.e. 加入一個元素後秩 不會降底,最多增加 1
- 5. $\forall Y \subseteq E, \exists X \subseteq Y, r(X) = |X| = r(Y)$ 一些等價的性質:
- 1. 對於所有 $X \subseteq E$, X 的極大獨立子集都有相同的大小
- 2. 對於 $B_1, B_2 \in \mathcal{B} \land B_1 \neq B_2$,對於所有 $e_1 \in B_1 \setminus B_2$,存在 $e_2 \in B_2 \setminus B_1$ 使 得 $(B_1 \setminus \{e_1\}) \cup \{e_2\} \in \mathcal{B}$
- 3. 對於 $X,Y \in \mathcal{I}$ 且 |X| < |Y|,存在 $e \in Y \setminus X$ 使得 $X \cup \{e\} \in \mathcal{B}$
- 4. 如果 $r(X \cup \{e_1\}) = r(X \cup \{e_2\}) = r(X)$,則 $r(X \cup \{e_1, e_2\}) = r(X)$ 。如果 $r(X \cup \{e\}) = r(X)$ 對於所有 $e \in E'$ 都成立,則 $r(X \cup E') = r(X)$ 。

```
Data: 兩個擬陣 M_1 = (E, \mathcal{I}_1) 以及 M_2 = (E, \mathcal{I}_2)
Result: I 為最大的位於 \mathcal{I}_1 \cap \mathcal{I}_2 中的獨立集
I \leftarrow \emptyset
X_1 \leftarrow \{e \in E \setminus I \mid I \cup \{e\} \in \mathcal{I}_1\}
X_2 \leftarrow \{e \in E \setminus I \mid I \cup \{e\} \in \mathcal{I}_2\}
while X_1 \neq \emptyset \coprod X_2 \neq \emptyset do
     if e \in X_1 \cap X_2 then
           I \leftarrow I \cup \{e\}
           構造交換圖 \mathcal{D}_{M_1,M_2}(I) 在交換圖上找到一條 X_1 到 X_2 且沒有捷徑的路徑 P
           I \leftarrow I \triangle P
     end if
     X_1 \leftarrow \{e \in E \setminus I \mid I \cup \{e\} \in \mathcal{I}_1\}
     X_2 \leftarrow \{e \in E \setminus I \mid I \cup \{e\} \in \mathcal{I}_2\}
```

Polynomial

\mathbf{FWHT} 8.1

end while

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
    for (int i = 0; i < n; i += L)
      for (int j = i; j < i + (L >> 1); ++j)
        a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];</pre>
void subset_convolution(int *a, int *b, int *c, int L) {
  // c_k = \sum_{i = 0} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
        h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i) fwt(h[i], n, -1);</pre>
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
8.2 FFT
// Errichto: FFT for double works when the result < 1e15,
    and < 1e18 with long double
using val_t = complex<double>;
template<int MAXN>
struct FFT {
  const double PI = acos(-1);
  val_t w[MAXN];
  FFT() {
    for (int i = 0; i < MAXN; ++i) {</pre>
      double arg = 2 * PI * i / MAXN;
      w[i] = val_t(cos(arg), sin(arg));
```

```
void bitrev(vector<val_t> &a, int n) //same as NTT
  void trans(vector<val_t> &a, int n, bool inv = false) {
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + dl; ++j, x += dx) {
          val_t + mp = a[j + dl] * (inv ? conj(w[x]) : w[x])
          a[j + dl] = a[j] - tmp;
          a[j] += tmp;
      }
    if (inv) {
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
  //multiplying two polynomials A * B:
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
  //A[i] *= B[i], fft.trans(A, siz, 1);
};
```

8.3 NTT

```
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
// only works when sz(A) + sz(B) - 1 <= MAXN
```

```
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
struct NTT {
  11 w[MAXN];
 11 mpow(ll a, ll n);
 11 minv(ll a) { return mpow(a, P - 2); }
   ll dw = mpow(RT, (P - 1) / MAXN);
   w[0] = 1;
   for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw % P
 void bitrev(vector<ll> &a, int n) {
   int i = 0;
    for (int j = 1; j < n - 1; ++j) {
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
   }
 void operator()(vector<ll> &a, int n, bool inv = false) {
       //0 <= a[i] < P
   bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
        for (int j = i, x = 0; j < i + dl; ++j, x += dx) {
          ll tmp = a[j + dl] * w[x] % P;
          if ((a[j + dl] = a[j] - tmp) < 0) a[j + dl] += P;
          if ((a[j] += tmp) >= P) a[j] -= P;
        }
     }
    if (inv) {
      reverse(a.begin()+1, a.begin()+n);
      11 invn = minv(n);
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn % P;</pre>
 }
```

8.4 Polynomial Operation

```
// Copy from 8BQube
#define fi(s, n) for (int i = (int)(s); i < (int)(n); ++i)
template<int MAXN, 11 P, 11 RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
 using vector<11>::vector;
 static inline NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); } // n() >= 1
 Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
 Poly& irev() { return reverse(data(), data() + n()), *
      this; }
 Poly& isz(int m) { return resize(m), *this; }
 Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if (((*this)[i] += rhs[i]) >= P) (*this)[i]
        -= P;
    return *this;
 Poly& imul(ll k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
    return *this;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    assert(m <= MAXN);</pre>
    Poly X(*this, m), Y(rhs, m);
    ntt(X, m), ntt(Y, m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X, m, true);
    return X.isz(n() + rhs.n() - 1);
 Poly Inv() const { // (*this)[0] != 0, 1e5/95ms, 2*sz<=
    if (n() == 1) return {ntt.minv((*this)[0])};
```

```
int m = 1;
  while (m < n() * 2) m <<= 1;</pre>
  assert(m <= MAXN);</pre>
  Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
  Poly Y(*this, m);
  ntt(Xi, m), ntt(Y, m);
  fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  ntt(Xi, m, true);
 return Xi.isz(n());
Poly& shift_inplace(const ll &c) { // 2 * sz <= MAXN
  int n = this->n();
  vector<ll> fc(n), ifc(n);
  fc[0] = ifc[0] = 1;
  for (int i = 1; i < n; i++){
    fc[i] = fc[i-1] * i % P;
    ifc[i] = ntt.minv(fc[i]);
  for (int i = 0; i < n; i++) (*this)[i] = (*this)[i] *
      fc[i] % P;
  Poly g(n);
  11 cp = 1;
  for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P, cp
      = cp * c % P;
  *this = (*this).irev().Mul(g).isz(n).irev();
  for (int i = 0; i < n; i++) (*this)[i] = (*this)[i] *
      ifc[i] % P;
  return *this;
Poly shift(const 11 &c) const { return Poly(*this).
    shift_inplace(c); }
Poly \_Sqrt() const { // Jacobi((*this)[0], P) = 1
  if (n() == 1) return {QuadraticResidue((*this)[0], P)};
  Poly X = Poly(*this, (n() + 1) / 2)._Sqrt().isz(n());
  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
Poly Sqrt() const { // 2 * sz <= MAXN
  Poly a;
  bool has = 0;
  for(int i = 0; i < n(); i++){</pre>
    if((*this)[i]) has = 1;
    if(has) a.push_back((*this)[i]);
  if(!has) return *this;
  if( (n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
    return Poly();
  a=a.isz((n() + a.n()) / 2)._Sqrt();
  int sz = a.n();
  a.isz(n());
  rotate(a.begin(), a.begin() + sz, a.end());
pair<Poly, Poly> DivMod(const Poly &rhs) const { // (rhs
    .)back() != 0
  if (n() < rhs.n()) return {{0}, *this};</pre>
  const int m = n() - rhs.n() + 1;
  Poly X(rhs); X.irev().isz(m);
  Poly Y(*this); Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
  X = rhs.Mul(Q), Y = *this;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
  fi(0, ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n()) ret[i + 1] = ntt.minv(i + 1) * (*this)[i] %
  return ret;
```

```
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const vector<ll> &x, const vector<Poly>
    &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i]).
      second:
  down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
      _tmul(m, *this);
  fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() - 1,
      down[i / 2]);
  vector<11> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].Mul(
     up[i * 2 + 1]);
  return up;
vector<ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate(const vector<11> &x, const vector
    &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
  vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m - 1; i > 0; --i) down[i] = down[i * 2].
      Mul(up[i * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2 + 1])
      2]));
  return down[1];
Poly Ln() const \{ // (*this)[0] == 1, 2*sz <= MAXN \}
  return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { // (*this)[0] == 0,2*sz<=MAXN
  if (n() == 1) return {1};
  Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
  Poly Y = X.Ln(); Y[0] = P - 1;
  fi(0, n()) if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] +=
  return X.Mul(Y).isz(n());
// M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(ll k) const { // 2*sz<=MAXN
  int nz = 0;
  while (nz < n() && !(*this)[nz]) ++nz;</pre>
  if (nz * min(k, (11)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly {1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const ll c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul(k % P).Exp().imul(c).irev().isz(n())
      .irev();
static 11 LinearRecursion(const vector<11> &a, const
    vector<ll> &coef, ll n) { // a_n = \sum a_n(n-j)
  const int k = (int)a.size();
  assert((int)coef.size() == k + 1);
  Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\};
  fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
  C[k] = 1;
  while (n)
    if (n % 2) W = W.Mul(M).DivMod(C).second;
    n /= 2, M = M.Mul(M).DivMod(C).second;
```

```
11 \text{ ret} = 0;
    fi(0, k) ret = (ret + W[i] * a[i]) % P;
    return ret;
#undef fi
using Poly_t = Poly<1 << 20, 998244353, 3>;
// template<> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

Generating Function 8.5

8.5.1 Ordinary Generating Function

- C(x) = A(rx): $c_n = r^n a_n$ 的一般生成函數。 • C(x) = A(x) + B(x): $c_n = a_n + b_n$ 的一般生成函數。
- C(x) = A(x)B(x): $c_n = \sum_{i=0}^n a_i b_{n-i}$ 的一般生成函數。 $C(x) = A(x)^k$: $c_n = \sum_{i_1+i_2+\dots+i_k=n} a_{i_1}a_{i_2}\dots a_{i_k}$ 的一般生成函數。
- C(x) = xA(x)': $c_n = na_n$ 的一般生成函數。
- $C(x) = \frac{A(x)}{1-x}$: $c_n = \sum_{i=0}^{n} a_i$ 的一般生成函數。
- $C(x) = A(1) + x \frac{A(1) A(x)}{1 x}$: $c_n = \sum_{i=-n}^{\infty} a_i$ 的一般生成函數。

- 常用展開式 $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$, ${a \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

• 卡特蘭數: $f(x) = \frac{1-\sqrt{1-4x}}{2x}$

8.5.2 Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)^k$: $\sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_i a_{i_2} \dots a_{i_k}$ 的指數生成函數
- $\hat{C}(x) = \exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

8.6Bostan Mori

```
NTT<262144, 998244353, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log k)
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k) {
  int d = max((int)P.size(), (int)Q.size() - 1);
  P.resize(d, 0);
  Q.resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) + 1)));
  while(k) {
    vector<ll> Qneg(sz);
    for(int i = 0; i < (int)Q.size(); i++){</pre>
      Qneg[i] = Q[i] * ((i & 1) ? -1 : 1);
      if(Qneg[i] < 0) Qneg[i] += mod;</pre>
    ntt(Qneg, sz, false);
    vector<ll> U(sz), V(sz);
    for(int i = 0; i < (int)P.size(); i++)</pre>
      U[i] = P[i];
    for(int i = 0; i < (int)Q.size(); i++)</pre>
      V[i] = Q[i];
    ntt(U, sz, false);
    ntt(V, sz, false);
    for(int i = 0; i < sz; i++)
U[i] = U[i] * Qneg[i] % mod;</pre>
```

```
for(int i = 0; i < sz; i++)
    V[i] = V[i] * Qneg[i] % mod;
ntt(U, sz, true);
ntt(V, sz, true);

for(int i = k & 1; i <= 2 * d - 1; i += 2)
    P[i >> 1] = U[i];
    for(int i = 0; i <= 2 * d; i += 2)
    Q[i >> 1] = V[i];
    k >>= 1;
}
return P[0] * ntt.minv(Q[0]) % mod;
}
```

9 String

9.1 KMP Algorithm

```
// 0-based
// fail[i] = max k < i s.t. s[0..k] = s[i-k..i]
vector<int> kmp_build_fail(const string &s){
  int n = SZ(s);
  vector<int> fail(n, -1);
  int cur = -1;
  for(int i = 1; i < n; i++){</pre>
    while(cur != -1 && s[cur + 1] != s[i])
      cur = fail[cur];
    if(s[cur + 1] == s[i])
      cur++;
    fail[i] = cur;
  return fail;
void kmp_match(const string &s, const vector<int> &fail,
    const string &t){
  int cur = -1;
  int n = SZ(s), m = SZ(t);
  for(int i = 0; i < m; i++){</pre>
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t[i])
      cur = fail[cur];
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
      cur++:
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
}
```

9.2 Manacher Algorithm

9.3 Lyndon Factorization

```
// partition s = w[0] + w[1] + ... + w[k-1], // w[0] >= w[1] >= ... >= w[k-1]
```

```
// each w[i] strictly smaller than all its suffix
void duval(const string &s, vector<pii> &w) {
    for (int n = (int)s.size(), i = 0, j, k; i < n; ) {
        for (j = i + 1, k = i; j < n && s[k] <= s[j]; j++)
            k = (s[k] < s[j] ? i : k + 1);
            // if (i < n / 2 && j >= n / 2) {
            // for min cyclic shift, call duval(s + s)
            // then here s.substr(i, n / 2) is min cyclic shift
            // }
        for (; i <= k; i += j - k)
            w.pb(pii(i, j - k)); // s.substr(l, len)
    }
}</pre>
```

9.4 Suffix Array

```
struct SuffixArray {
  vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i] and
      sa[i-1]
                              // sa[0] = s.size()
                               // character should be 1-based
  SuffixArray(string& s, int lim=256) { // or basic_string
    int n = s.size() + 1, k = 0, a, b;
    vector<int> x(n, 0), y(n), ws(max(n, lim));
    rank.assign(n, 0);
    for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
    sa = lcp = y, iota(sa.begin(), sa.end(), 0);
    for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = 0
        p) {
      p = j, iota(y.begin(), y.end(), n - j);
      for (int i = 0; i < n; i++)</pre>
        if (sa[i] >= j) y[p++] = sa[i] - j;
      for (int &i : ws) i = 0;
      for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
      for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];</pre>
      for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
      swap(x, y), p = 1, x[sa[0]] = 0;
      for(int i = 1; i < n; i++){</pre>
        a = sa[i - 1], b = sa[i];
        x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ? p -
             1 : p++;
    for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
    for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
      for (k && k--, j = sa[rank[i] - 1];
          s[i + k] == s[j + k]; k++);
 }
};
```

9.5 Suffix Automaton

```
struct exSAM {
 const int CNUM = 26;
  // len: maxlength, link: fail link
 // LenSorted: topo order, cnt: occur
  vector<int> len, link, lenSorted, cnt;
 vector<vector<int>> next;
  int total = 0;
 int newnode() {
   return total++;
 void init(int n) { // total number of characters
    len.assign(2 * n, 0); link.assign(2 * n, 0);
    lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
    next.assign(2 * n, vector<int>(CNUM));
   newnode(), link[0] = -1;
 int insertSAM(int last, int c) {
   // not exSAM: cur = newnode(), p = last
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
```

```
next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len[p] + 1 == len[q]) return link[cur] = q, cur;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)
      next[clone][i] = len[next[q][i]] ? next[q][i] : 0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
 void insert(const string &s) {
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
  void build() {
    queue<int> q;
    q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    }
    vector<int> lc(total);
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[i]]]</pre>
 void solve() {
    for (int i = total - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
9.6 Z-value Algorithm
```

```
// z[i] = max k s.t. s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z_function(const string &s){
 int n = s.size();
  vector<int> z(n);
  for(int i = 1, l = 0, r = 0; i < n; i++){
    if(i \leftarrow r) z[i] = min(r - i + 1, z[i - 1]);
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
     z[i]++;
    if(i + z[i] - 1 > r)
      l = i, r = i + z[i] - 1;
 return z;
```

9.7Main Lorentz

```
struct Rep{ int minl, maxl, len; };
vector<Rep> rep; // 0-base
// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)
void main_lorentz(const string &s, int sft = 0) {
 const int n = s.size();
  if (n == 1) return;
 const int nu = n / 2, nv = n - nu;
 const string u = s.substr(0, nu), v = s.substr(nu),
        ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend());
 main_lorentz(u, sft), main_lorentz(v, sft + nu);
```

```
const auto z1 = z_function(ru), z2 = z_function(v + '#' +
             z3 = z function(ru + '#' + rv), z4 =
                  z_function(v);
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0; };</pre>
  auto add_rep = [&](bool left, int c, int l, int k1, int
      k2) {
    const int L = max(1, 1 - k2), R = min(1 - left, k1);
    if (L > R) return;
    if (left) rep.emplace_back(Rep({sft + c - R, sft + c -
        L, 1}));
    else rep.emplace_back(Rep({sft + c - R - l + 1, sft + c
         -L-l+1, 1\}));
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int 1, k1, k2;
    if (cntr < nu) {
      1 = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
}
```

AC Automaton 9.8

```
const int SIGMA = 26;
struct AC_Automaton {
  // child: trie, next: automaton
  vector<vector<int>> child, next;
  vector<int> fail, cnt, ord;
  int total = 0;
  int newnode() {
    return total++;
  void init(int len) { // len >= 1 + total len
    child.assign(len, vector<int>(26, -1));
    next.assign(len, vector<int>(26, -1));
    fail.assign(len, -1); cnt.assign(len, 0);
    ord.clear():
    newnode();
  int input(string &s) {
    int cur = 0;
    for (char c : s) {
      if (child[cur][c - 'A'] == -1)
    child[cur][c - 'A'] = newnode();
      cur = child[cur][c - 'A'];
    return cur; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(0), fail[0] = -1;
    while(!q.empty()) {
      int R = q.front();
      q.pop(); ord.pb(R);
for (int i = 0; i < SIGMA; i++)</pre>
        if (child[R][i] != -1) {
           int X = next[R][i] = child[R][i], Z = fail[R];
          while (Z != -1 && child[Z][i] == -1)
             Z = fail[Z];
          fail[X] = Z != -1 ? child[Z][i] : 0;
          q.push(X);
        else next[R][i] = R ? next[fail[R]][i] : 0;
  void solve() {
```

```
for (int i : ord | views::reverse)
 cnt[fail[i]] += cnt[i];
```

Palindrome Automaton

```
struct PalindromicTree {
  struct node {
   int nxt[26], fail, len; // num = depth of fail link
    int cnt, num; // cnt = occur, num = #pal_suffix of this
   node(int 1 = 0) : nxt{}, fail(0), len(1), cnt(0), num(0) {}
 };
 vector<node> st; vector<int> s; int last, n;
 void init() {
    st.clear(); s.clear(); last = 1; n = 0;
   st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
 int getFail(int x) {
   while (s[n - st[x].len - 1] != s[n]) x = st[x].fail;
   return x;
 void add(int c) {
    s.pb(c -= 'a'); ++n;
   int cur = getFail(last);
   if (!st[cur].nxt[c]) {
     int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
      st[now].num = st[st[now].fail].num + 1;
   last = st[cur].nxt[c]; ++st[last].cnt;
 void dpcnt() {
    for(int i = SZ(st) - 1; i >= 0; i--){
     auto nd = st[i];
      st[nd.fail].cnt += nd.cnt;
 int size() { return (int)st.size() - 2; }
```

Formula 10

Recurrences

If $a_n = c_1 a_{n-1} + \cdots + c_k a_{n-k}$, and r_1, \dots, r_k are distinct roots of $x^k + c_1 x^{k-1} + \cdots + c_k a_{n-k}$ $\cdots + c_k$, there are d_1, \ldots, d_k s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1n + d_2)r^n$.

10.2Geometry

Rotation Matrix

$$\begin{pmatrix}
\cos\theta & -\sin\theta \\
\sin\theta & \cos\theta
\end{pmatrix}$$

- rotate 90°: $(x,y) \rightarrow (-y,x)$
- rotate -90° : $(x,y) \rightarrow (y,-x)$

10.2.2 Triangles

Side lengths:
$$a,b,c$$

Semiperimeter: $p = \frac{a+b+c}{2}$
Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Area:
$$A = \sqrt{p(p-a)(p-b)(p-a)}$$

Circumradius: $R = \frac{abc}{c}$

Inradius:
$$r = \frac{A}{p}$$

Length of median (divides triangle into two equal-area triangles): $m_a =$

Length of bisector (divides angles in two): $s_a = \sqrt{bc \left(1 - \left(\frac{a}{b+c}\right)^2\right)}$

Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$ a b cLaw of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$ Incenter:

 $P_1 = (x_1, y_1), P_2 = (x_2, y_2), P_3 = (x_3, y_3)$ $s_1 = \overline{P_2 P_3}, s_2 = \overline{P_1 P_3}, s_3 = \overline{P_1 P_2}$ $s_1 P_1 + s_2 P_2 + s_3 P_3$

 $s_1 + s_2 + s_3$

Circumcenter: $P_0 = (0,0), P_1 = (x_1, y_1), P_2 = (x_2, y_2)$ $x_c = \frac{1}{2} \times \frac{y_2(x_1^2 + y_1^2) - y_1(x_2^2 + y_2^2)}{y_1^2 + y_2^2}$

 $y_c - \frac{1}{2} \wedge -x_1y_2 + x_2y_1$ Check if (x_0, y_0) is in the circumcircle:

$$\begin{vmatrix} x_1-x_0 & y_1-y_0 & (x_1^2+y_1^2)-(x_0^2+y_0^2) \\ x_2-x_0 & y_2-y_0 & (x_2^2+y_2^2)-(x_0^2+y_0^2) \\ x_3-x_0 & y_3-y_0 & (x_3^2+y_3^2)-(x_0^2+y_0^2) \end{vmatrix}$$

0: on edge, > 0: inside, < 0: outside

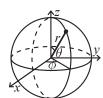
10.2.3 Quadrilaterals

With side lengths a, b, c, d, diagonals e, f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}.$

10.2.4 Spherical coordinates



$$x = r \sin \theta \cos \phi$$

$$y = r \sin \theta \sin \phi$$

$$z = r \cos \theta$$

$$r = \sqrt{x^2 + y^2 + z^2}$$

$$\theta = a\cos(z/\sqrt{x^2 + y^2 + z^2})$$

$$\phi = a\tan(y, x)$$

10.2.5 Green's Theorem

$$\iint_{D} \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy = \oint_{L^{+}} (Pdx + Qdy)$$

$$Area = \frac{1}{2} \oint_{I} x \ dy - y \ dx$$

Circular sector:

$$x = x_0 + r \cos \theta$$

$$y = y_0 + r \sin \theta$$

$$A = r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \ d\theta$$

$$= r(r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta}$$

10.2.6 Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1, p_2, p_3 are collinear $\iff p_1^*, p_2^*, p_3^*$ intersect at a point
- p lies above $l \iff l^*$ lies above p^*
- lower convex hull \leftrightarrow upper envelope

10.3 Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x}) \qquad \cosh x = \frac{1}{2}(e^x + e^{-x})$$

$$\sin n\pi = 0 \qquad \cos n\pi = (-1)^n$$

$$\sin(\alpha + \beta) = \sin \alpha \cos \beta + \cos \alpha \sin \beta$$

$$\cos(\alpha + \beta) = \cos \alpha \cos \beta - \sin \alpha \sin \beta$$

$$\sin(2\alpha) = 2\cos \alpha \sin \alpha$$

$$\cos(2\alpha) = \cos^2 \alpha - \sin^2 \alpha$$

$$= 2\cos^2 \alpha - 1$$

$$= 1 - 2\sin^2 \alpha$$

$$\tan(\alpha + \beta) = \frac{\tan \alpha + \tan \beta}{1 - \tan \alpha \tan \beta}$$

$$\sin \alpha + \sin \beta = 2\sin \frac{\alpha + \beta}{2}\cos \frac{\alpha - \beta}{2}$$

$$\cos \alpha + \cos \beta = 2\cos \frac{\alpha + \beta}{2}\cos \frac{\alpha - \beta}{2}$$

$$\sin \alpha \sin \beta = \frac{1}{2}(\cos(\alpha - \beta) - \cos(\alpha + \beta))$$

$$\sin \alpha \cos \beta = \frac{1}{2}(\sin(\alpha + \beta) + \sin(\alpha - \beta))$$

$$\cos \alpha \sin \beta = \frac{1}{2}(\sin(\alpha + \beta) - \sin(\alpha - \beta))$$

$$\cos \alpha \cos \beta = \frac{1}{2}(\cos(\alpha - \beta) + \cos(\alpha + \beta))$$

$$(V + W)\tan(\alpha - \beta)/2 = (V - W)\tan(\alpha + \beta)/2$$
where V, W are lengths of sides opposite angles α, β .
$$a \cos x + b \sin x = r \cos(x - \phi)$$

$$a \sin x + b \cos x = r \sin(x + \phi)$$
where $r = \sqrt{a^2 + b^2}, \phi = \tan 2(b, a)$.

10.4 Derivatives/Integrals

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx} \tan x = 1 + \tan^2 x \qquad \frac{d}{dx} \arctan x = \frac{1}{1+x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int xe^{ax} = \frac{e^{ax}}{a^2} (ax - 1)$$

$$\int \sin^2(x) = \frac{x}{2} - \frac{1}{4} \sin 2x \qquad \int \sin^3 x = \frac{1}{12} \cos 3x - \frac{3}{4} \cos x$$

$$\int \cos^2(x) = \frac{x}{2} + \frac{1}{4} \sin 2x \qquad \int \cos^3 x = \frac{1}{12} \sin 3x + \frac{3}{4} \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int xe^x = e^x(x - 1) \qquad \int x^2 e^x = e^x(x^2 - 2x + 2)$$

$$\int x^2 \sin x = 2x \sin x - (x^2 - 2) \cos x$$

$$\int x^2 \cos x = 2x \cos x + (x^2 - 2) \sin x$$

$$\int e^x \sin x = \frac{1}{2} e^x (\sin x - \cos x)$$

$$\int e^x \cos x = \frac{1}{2} e^x (\sin x + \cos x)$$

$$\int x e^x \sin x = \frac{1}{2} e^x (x \sin x - x \cos x + \cos x)$$

$$\int x e^x \cos x = \frac{1}{2} e^x (x \sin x + x \cos x - \sin x)$$

10.5 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

10.6 Series

$$\begin{split} e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, \, (-\infty < x < \infty) \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, \, (-1 < x \le 1) \\ \sqrt{1+x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, \, (-1 \le x \le 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \, (-\infty < x < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \, (-\infty < x < \infty) \end{split}$$