using pll = pair<ll, ll>;

void debug(){cerr << "\n";}</pre> template < class T, class ... U>

#ifdef zisk

void debug(T a, U ... b){cerr << a << " ", debug(b...)</pre> **Contents** 5.11 Minimum Steiner Tree . . . 14 5.12 Count Cycles 15 Basic 6 Math 1.1 Default Code 6.1 Extended Euclidean Algo-1.2 .vimrc Fast IO rithm 15 13 1.4 Random 6.2 Floor & Ceil 15 1.5 PBDS Tree 6.3 Legendre 15 Pragma 6.5 Simplex Construction . . . 16 6.6 DiscreteLog 16 2 Data Structure 2.1 Heavy-Light Decomposition 6.7 Miller Rabin & Pollard Rho . 16 2.2 Link Cut Tree 6.8 XOR Basis 16 2.3 6.9 Linear Equation 17 2.4 KD Tree 6.10 Chinese Remainder Theorem 17 2.5 Leftist Tree2.6 Convex 1D/1D2.7 Dynamic Convex Hull . . . 6.11 Sqrt Decomposition 17 3 Flow & Matching 7 Polynomial 3.1 Dinic . . 7.1 FWHT 17 3.2 Bounded Flow 7.2 FFT 18 MCMF 3.3 7.3 NTT 18 Min Cost Circulation Gomory Hu 7.4 Polynomial Operation . . . 18 7.5 Generating Function 20 3.6 Stoer Wagner Algorithm . . 3.7 Bipartite Matching Ordinary Generating Func-3.8 Kuhn Munkres Algorithm . tion 20 Exponential Generating 6 Max Simple Graph Matching 3.9 3.10 Flow Model Function 20 7.6 Bostan Mori 20 Geometry 4.1 Geometry Template 8 String Polar Angle Comparator . . Minkowski Sum 8.1 KMP Algorithm 21 Intersection of Circle and 8.2 Manacher Algorithm 21 8.3 Lyndon Factorization 21 8.4 Suffix Array 21 Tangent Line of Circles . . . 4.6 8.5 Suffix Automaton 21 4.7 Intersection of Line and 8.6 Z-value Algorithm 22 Convex Polygon 4.8 Intersection of Line and 8.7 Main Lorentz 22 Circle 8.8 AC Automaton 22 4.9 Point in Circle 8.9 Palindrome Automaton . . 22 4.10 Point in Convex 4.11 Half Plane Intersection 4.12 Minimum Enclosing Circle . 9.1 Cyclic Ternary Search . . . 23 9.2 Matroid 23 4.15 Delaunay Triangulation . . 10 9.3 Simulate Annealing 23 4.16 Voronoi Diagram 10 9.4 Binary Search On Fraction . 23 4.17 Polygon Union 9.5 Min Plus Convolution . . . 23 4.18 Tangent Point to Convex Hull 11 9.6 SMAWK 23 4.19 Heart 11 4.20 Rotating Sweep Line 11 4.23 Calculate Points in Triangle 12 Rotation Matrix 24 Triangles 24 Graph Quadrilaterals 24 5.2 SCC 12 5.3 2-SAT 12 5.4 Dominator Tree 12 Spherical coordinates . . . 24 Green's Theorem 24 Point-Line Duality 24 10.2 Trigonometry 24 5.6 Fast DMST 13 10.3 Calculus 24 10.4 Sum & Series 24 5.9 Number of Maximal Clique 14 5.10 Minimum Mean Cycle . . . 14 10.5 Misc 25 10.6 Number 25 Basic 1.1 Default Code [2d2229] #include <bits/stdc++.h> using namespace std; #define iter(v) v.begin(),v.end() #define SZ(v) int(v.size()) #define pb emplace_back #define ff first #define ss second using 11 = long long; using pii = pair<int, int>;

```
template < class T> void pary(T 1, T r){
  while (1 != r) cerr << *1 << "
  cerr << "\n";</pre>
#else
#define debug(...) void()
#define pary(...) void()
#endif
template < class A, class B>
ostream& operator<<(ostream& o, pair<A,B> p) {    return o << '(' << p.ff << ',' << p.ss <<
                                  << p.ss << ')'; }
int main(){
  ios_base::sync_with_stdio(false); cin.tie(0);
1.2 .vimrc [b4816d]
se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a
map <F9> :w<bar>!g++ "%" -o %:r -std=c++17 -Wall -
    Wextra -Wshadow -O2 -Dzisk -g -fsanitize=address,
    undefined<CR>
map <F8> :!./%:r<CR>
inoremap {<CR> {<CR>}<ESC>ko
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space
    :]' \| md5sum \| cut -c-6
inoremap fj <ESC>
vnoremap fj <ESC>
" -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG
1.3 Fast IO [4f6f0e]
char readchar() {
  const int N = 1<<20;</pre>
  static char buf[N];
  static char *p = buf , *end = buf;
  if(p == end) {
    if((end = buf + fread(buf , 1 , N , stdin)) == buf)
         return EOF;
    p = buf;
  return *p++;
}
const int buf size = 524288;
struct Writer {
  char buf[buf_size]; int size = 0, ret;
  void flush() { ret = write(1, buf, size); size = 0; }
  void _flush(int sz) { if (sz + size > buf_size) flush
      (); }
  void write_char(char c) { _flush(1); buf[size++] = c;
  void write_int(int x) {
    const int len = 20;
     _flush(len);    <mark>int</mark> ptr = 0;
    if (x < 0) buf[size++] = '-', x = -x;</pre>
    if (x == 0) buf[size + (ptr++)] = '0';
    else for (; x; x /= 10) buf[size + (ptr++)] = '0' +
         x % 10;
    reverse(buf + size, buf + size + ptr);
    size += ptr:
}; // remember to call flush
1.4 Random [4cf9ed]
mt19937 rng(chrono::system clock::now().
    time_since_epoch().count());
1.5 PBDS Tree [9e57e3]
#include <bits/extc++.h>
using namespace __gnu_pbds;
using Tree = tree<int, null_type, less<>, rb_tree_tag,
    tree_order_statistics_node_update>;
// .find_by_order(x)
// .order_of_key(x)
```

ofstream o; string c = "red";

public: // SVG svg("test.svg", 0, 0, 100, 100)

1.6 Pragma [6006f6]

```
#pragma GCC optimize("Ofast,no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popent,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)

1.7 SVG Writer [7adcc8]

class SVG {
   void p(string_view s) { o << s; }
   void p(string_view s, auto v, auto... vs) {
     auto i = s.find('$');
     o << s.substr(0, i) << v, p(s.substr(i + 1), vs...)</pre>
```

SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f) { p("<svg xmlns='http://www.w3.org/2000/svg' " "viewBox='\$ \$ \$'>\n" "<style>*{stroke-width:0.5%;}</style>\n", x1, -y2, x2 - x1, y2 - y1); } ~SVG() { p("</svg>\n"); } void color(string nc) { c = nc; } void line(auto x1, auto y1, auto x2, auto y2) { p("kine x1='\$' y1='\$' x2='\$' y2='\$' stroke='\$'/>\n", x1, -y1, x2, -y2, c); } void circle(auto x, auto y, auto r) { p("<circle cx='\$' cy='\$' r='\$' stroke='\$' " "fill='none'/>\n", x, -y, r, c); } void text(auto x, auto y, strings, int w = 12) { "fill='none'/>\n", x, -y, r, c); } void text(auto x, auto y, strings, int w = 12) { "fill='none'/>\n", x, -y, r, c); }

p("<text x='\$' y='\$' font-size='\$px'>\$</text>\n",

2 Data Structure

};

x, -y, w, s); }

2.1 Heavy-Light Decomposition [f2dbca]

```
struct HLD{ // 1-based
  int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
  explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1),
  ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
   }
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  void build(){ dfs(1, 1), cut(1, 1); }
int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // Lca
};
```

2.2 Link Cut Tree [cf4f34]

```
// 1-based
// == 43515a ==
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node() : pa{0}, ch{0, 0}, rev{false}, size{1}, v{},
      sum\{\}, rsum\{\}, sv\{\}, sub\{\}, vir\{\} \{\}
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is_rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false:
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
    cur.rev ^= 1;
// == f05d4f ==
  void rotate(int u) {
    int f = cur.pa, g = o[f].pa, l = is_rch(u);
    if (cur.ch[1 ^ 1]) o[cur.ch[1 ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  void splay(int u) {
    vector<int> stk = {u};
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is_root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    up(u);
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
      cur.vir = cur.vir + o[rc].sub - o[last].sub;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0:
    for (access(u); u; u = lc) down(la = u);
    return la;
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
// == a238c2 ==
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); }
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
```

```
return find_root(u) == find_root(v); }
  void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
#undef rc
};
2.3 Treap [2ac37e]
mt19937 rng(880301);
// == fb4359 ==
struct node {
  11 data; int sz;
  node *1, *r;
  node(11 k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    sz = 1;
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
node pool[1000010]; int pool_cnt = 0;
node *newnode(11 k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
  return b->down(), b->l = merge(a, b->l), b->up(), b;
}
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
  if (!o) return a = b = 0, void();
  o->down();
  if (o->data <= k)</pre>
    a = o, split(o->r, a->r, b, k), a->up();
  else b = o, split(o->1, a, b->1, k), b->up();
// a: size k, b: size n - k
void split2(node *o, node *&a, node *&b, int k) {
  if (sz(o) <= k) return a = o, b = 0, void();</pre>
  o->down();
  if (sz(o->1) + 1 <= k)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
}
// == e9f4d8 ==
node *kth(node *o, ll k) { // 1-based
  if (k <= sz(o->1)) return kth(o->1, k);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
int Rank(node *o, 11 key) { // num of key < key</pre>
  if (!o) return 0;
  if (o->data < key)</pre>
    return sz(o->1) + 1 + Rank(o->r, key);
  else return Rank(o->1, key);
bool erase(node *&o, ll k) {
  if (!o) return 0;
  if (o->data == k) {
    node *t = o;
    o->down(), o = merge(o->1, o->r);
    return 1;
  node *&t = k < o->data ? o->l : o->r;
  return erase(t, k) ? o->up(), 1 : 0;
void insert(node *&o, ll k) {
  node *a, *b;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
```

```
tuple<node*, node*, node*> interval(node *&o, int 1,
    int r) { // 1-based
  node *a, *b, *c; // b: [l, r]
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
  return make_tuple(a, b, c);
}
2.4 KD Tree [375ca2]
namespace kdt {
  int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
  point p[maxn];
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
    function < bool (const point &, const point &) > f =
      [dep](const point &a, const point &b) {
        if (dep & 1) return a.x < b.x;</pre>
        else return a.y < b.y;</pre>
      };
    int m = (1 + r) >> 1;
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
    yl[m] = yr[m] = p[m].y;
    lc[m] = build(1, m, dep + 1);
    if (~lc[m]) {
      x1[m] = min(x1[m], x1[1c[m]]);
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
      yr[m] = max(yr[m], yr[lc[m]]);
    rc[m] = build(m + 1, r, dep + 1);
    if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
      yr[m] = max(yr[m], yr[rc[m]]);
    return m;
  bool bound(const point &q, int o, long long d) {
    double ds = sqrt(d + 1.0);
    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
        q.y < yl[o] - ds || q.y > yr[o] + ds)
      return false;
    return true;
  long long dist(const point &a, const point &b) {
    return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
  void dfs(
      const point &q, long long &d, int o, int dep = 0)
    if (!bound(q, o, d)) return;
    long long cd = dist(p[o], q);
    if (cd != 0) d = min(d, cd);
    if ((dep & 1) && q.x < p[o].x ||</pre>
        !(dep & 1) && q.y < p[o].y) {
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    } else {
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    }
  void init(const vector<point> &v) {
    for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
    root = build(0, v.size());
  long long nearest(const point &q) {
    long long res = 1e18;
    dfs(q, res, root);
    return res:
} // namespace kdt
2.5 Leftist Tree [e91538]
struct node {
  11 v, data, sz, sum;
node *1, *r;
  node(ll k)
```

}:

```
: v(0), data(k), sz(1), l(0), r(0), sum(k) {}
11 sz(node *p) { return p ? p->sz : 0; }
11 V(node *p) { return p ? p->v : -1; }
11 sum(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
  if (!a || !b) return a ? a : b;
  if (a->data < b->data) swap(a, b);
  a->r = merge(a->r, b);
  if (V(a\rightarrow r) \rightarrow V(a\rightarrow l)) swap(a\rightarrow r, a\rightarrow l);
  a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  return a:
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->1, o->r);
  delete tmp;
}
2.6 Convex 1D/1D [a449dd]
```

```
template < class T>
struct DynamicHull {
  struct seg { int x, l, r; };
  T f; int C; deque<seg> dq; // range: 1~C
  explicit DynamicHull(T _f, int _C): f(_f), C(_C) {} // max t s.t. f(x, t) >= f(y, t), x < y, maintain max
  int intersect(int x, int y) {
    int 1 = 0, r = C + 1;
    while (1 + 1 < r) {
      int mid = (1 + r) / 2;
      if (f(x, mid) >= f(y, mid)) l = mid;
      else r = mid;
    }
    return 1;
  void push_back(int x) {
    for (int i; !dq.empty() &&
         (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
         ();
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back(). r + 1 \le C) dq.pb(seg({x, dq.back().}
        r + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
  }
};
```

Dynamic Convex Hull [b45ebc]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
      rhs.a; }
  bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  bool isect(iterator x, iterator y) {
    if (y == end()) \{ x \rightarrow p = kInf; return 0; \}
    if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf
    else x -> p = iceil(y -> b - x -> b, x -> a - y -> a);
    return x->p >= y->p;
  }
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) isect(x, y =
        erase(y));
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  11 query(ll x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
```

```
Flow & Matching
3
```

3.1 Dinic [801a71]

```
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
    ll to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
  vector<int> dis, ind;
  void init(int _n) {
    n = _n;
    g.assign(n, vector<edge>());
  void reset() {
    for (int i = 0; i < n; ++i)</pre>
      for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
    g[u].pb(edge{v, cap, 0, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
    //change g[v] to cap for undirected graphs
  bool bfs() {
    dis.assign(n, -1);
    queue<int> q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int cur = q.front(); q.pop();
      for (auto &e : g[cur])
        if (dis[e.to] == -1 && e.flow != e.cap) {
          q.push(e.to);
          dis[e.to] = dis[cur] + 1;
        }
      }
    }
    return dis[t] != -1;
  11 dfs(int u, ll cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df;
          g[e.to][e.rev].flow -= df;
          return df;
      }
    dis[u] = -1;
    return 0:
  11 maxflow(int _s, int _t) {
         _s; t = _t;
    11 \text{ flow} = 0, df;
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  }
};
3.2
      Bounded Flow [758826]
```

```
struct BoundedFlow : Dinic {
 vector<ll> tot;
 void init(int _n) {
   Dinic::init(_n + 2);
    tot.assign(n, 0);
 void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
```

```
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
                                                             struct MinCostCirculation { // 0-based, O(VE * ElogC)
                                                              struct edge {
  bool feasible() {
                                                                11 from, to, cap, fcap, flow, cost, rev;
    11 \text{ sum} = 0;
    int vs = n - 2, vt = n - 1;
                                                              int n;
    for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
        add_edge(vs, i, 0, tot[i]), sum += tot[i];
      else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
        g[vs].pop_back(), g[i].pop_back();
      else if(tot[i] < 0)</pre>
        g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
                                                                }:
    add_edge(_t, _s, 0, INF);
    if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    };
                                                                }
                                                              }
3.3 MCMF [671e14]
struct MCMF { // 0-base
  struct Edge {
   11 from, to, cap, flow, cost, rev;
  int n, s, t;
  vector<vector<Edge>> g;
  vector<Edge*> past;
  vector<ll> dis, up, pot;
                                                                  }
  explicit MCMF(int
                     _n): n(_n), g(n), past(n), dis(n),
      up(n), pot(n) {}
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(Edge{a, b, cap, 0, cost, SZ(g[b])});
    g[b].pb(Edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
  bool BellmanFord() {
    vector<bool> inq(n);
    fill(iter(dis), INF);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
                                                                }
        if (!inq[u]) inq[u] = 1, q.push(u);
                                                              }
      }
    };
                                                                n = n:
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
relax(e.to, d2, min(up[u], e.cap - e.flow), &e)
                                                              }
      }
                                                            };
    }
                                                            3.5
    return dis[t] != INF;
  pair<ll, 1l> solve(int _s, int _t, bool neg = true) {
    s = _s, t = _t; 11 flow = 0, cost = 0;
    if (neg) BellmanFord(), pot = dis;
    for (; BellmanFord(); pot = dis) {
      for (int i = 0; i < n; ++i)</pre>
        if (dis[i] != INF) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], g[e.to][e.rev].flow -= up[t];
                                                            }
      }
    return {flow, cost};
};
                                                              int n;
```

vector<edge*> past; vector<vector<edge>> g; vector<ll> dis; void BellmanFord(int s) { vector<int> inq(n); dis.assign(n, INF); queue<int> q; auto relax = [&](int u, ll d, edge *e) { **if** (dis[u] > d) { dis[u] = d, past[u] = e; **if** (!inq[u]) inq[u] = 1, q.push(u); relax(s, 0, 0); while (!q.empty()) { int u = q.front(); q.pop(), inq[u] = 0;for (auto &e : g[u]) if (e.cap > e.flow) relax(e.to, dis[u] + e.cost, &e); void try_edge(edge &cur) { if (cur.cap > cur.flow) return ++cur.cap, void(); BellmanFord(cur.to); if (dis[cur.from] + cur.cost < 0) {</pre> ++cur.flow, --g[cur.to][cur.rev].flow; for (int i = cur.from; past[i]; i = past[i]->from) { auto &e = *past[i]; ++e.flow, --g[e.to][e.rev].flow; ++cur.cap; void solve(int mxlg) { // mxlg >= log(max cap) for (int b = mxlg; b >= 0; --b) { for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) e.cap *= 2, e.flow *= 2; for (int i = 0; i < n; ++i)</pre> for (auto &e : g[i]) if (e.fcap >> b & 1) try_edge(e); void init(int _n) { past.assign(n, nullptr); g.assign(n, vector<edge>()); void add_edge(ll a, ll b, ll cap, ll cost) { $g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a$ == b)}): g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1}); **Gomory Hu** [82d968] void GomoryHu(Dinic &flow) { // 0-based int n = flow.n; vector<int> par(n); for (int i = 1; i < n; ++i) {</pre> flow.reset(); add_edge(i, par[i], flow.maxflow(i, par[i])); for (int j = i + 1; j < n; ++j)</pre> if (par[j] == par[i] && ~flow.dis[j]) par[j] = i; Stoer Wagner Algorithm [a9917b] struct StoerWagner { // 0-based, O(V^3) vector<int> vis, del; vector<ll> wei;

Min Cost Circulation [47cf18] 3.4

}:

```
vector<vector<ll>> edge;
  void init(int _n) {
   n = _n;
    del.assign(n, 0);
    edge.assign(n, vector<ll>(n));
  void add_edge(int u, int v, ll w) {
   edge[u][v] += w, edge[v][u] += w;
  void search(int &s, int &t) {
   vis.assign(n, 0); wei.assign(n, 0);
    s = t = -1;
    while (1) {
      11 mx = -1, cur = 0;
      for (int i = 0; i < n; ++i)</pre>
        if (!del[i] && !vis[i] && mx < wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)</pre>
        if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
   }
  ll solve() {
   11 ret = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
      for (int j = 0; j < n; ++j)</pre>
        edge[x][j] = (edge[j][x] += edge[y][j]);
    return ret:
 }
};
3.7
     Bipartite Matching [5bb9be]
// O(E sqrt(V)), O(E log V) for random sparse graphs
struct BipartiteMatching { // 0-based
  int nl, nr;
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!~my[e] || (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  }
  bool bfs() {
    int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < n1; ++i)</pre>
     if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
   }
    return ret;
  int matching() {
    int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
      cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
   n1 = _n1, nr = _nr;
    g.assign(nl, vector<int>());
```

```
Kuhn Munkres Algorithm [683e0a]
```

```
struct KM { // 0-based, maximum matching, O(V^3)
  int n, ql, qr;
  vector<vector<ll>> w;
  vector<ll> hl, hr, slk;
  vector<int> fl, fr, pre, qu, vl, vr;
  void init(int _n) {
    // -INF for perfect matching
    w.assign(n, vector<ll>(n, 0));
    pre.assign(n, 0);
    qu.assign(n, 0);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return (vr[qu[qr++] = fl[x]] = 1);
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
        0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (11 d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] -
               w[x][y])) {
            if (pre[x] = y, d) slk[x] = d;
            else if (!check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] \&\& d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (v1[x]) h1[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !s1k[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
         ; hr.assign(n, 0);
    for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
  3
};
      Max Simple Graph Matching [907d7c]
```

3.9

```
struct Matching { // 0-based, O(V^3)
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> g;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
   }
  void Blossom(int x, int y, int 1) {
    for (; Find(x) != 1; x = pre[y]) {
      pre[x] = y, y = match[x];
      if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
```

```
}
  bool Bfs(int r) {
    iota(iter(fa), 0); fill(iter(s), -1);
    q = queue<int>(); q.push(r); s[r] = 0;
    for (; !q.empty(); q.pop()) {
      for (int x = q.front(); int u : g[x])
        if (s[u] == -1) {
           if (pre[u] = x, s[u] = 1, match[u] == n) {
             for (int a = u, b = x, last;
    b != n; a = last, b = pre[a])
               last = match[b], match[b] = a, match[a] =
                    b;
             return true;
          }
           q.push(match[u]); s[match[u]] = 0;
        } else if (!s[u] && Find(u) != Find(x)) {
           int 1 = LCA(u, x);
           Blossom(x, u, 1); Blossom(u, x, 1);
    return false;
  Matching(\textbf{int} \_n) : n(\_n), fa(n + 1), s(n + 1), vis(n
  + 1), pre(n + 1, n), match(n + 1, n), g(n) {} void add_edge(int u, int v)
  { g[u].pb(v), g[v].pb(u); }
  int solve() {
    int ans = 0;
    for (int x = 0; x < n; ++x)
      if (match[x] == n) ans += Bfs(x);
    return ans;
  } // match[x] == n means not matched
};
```

3.10 Flow Model

- Maximum/Minimum flow with lower bound / Circulation problem
 - 1. Construct super source S and sink T.
- 2. For each edge (x,y,l,u), connect $x \to y$ with capacity u-l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
- 1. Consruct super source S and sink T
- 2. For each edge (x,y,c), connect x o y with (cost,cap)=(c,1) if c>0, otherwise connect y o x with (cost,cap)=(-c,1)
- 3. For each edge with c<0, sum these cost as K , then increase d(y) by 1. decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0,d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let ${\cal K}$ be the sum of all weights
- 3. Connect source $s \to v$, $v \in G$ with capacity K
- 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} u 2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
- 1. Let $w'(u,v)=w(u,v)-\mu(u)-\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.

- 2. Find the minimum weight matching M with w' . The answer is $\sum \mu(v) + w'(M)$.
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
 - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
- 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\begin{aligned} \min \sum_{uv} w_{uv} f_{uv} \\ -f_{uv} \geq -c_{uv} &\Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

4 Geometry

4.1 Geometry Template [86f0f1]

```
using ld = 11;
using pdd = pair<ld, ld>;
#define X first
#define Y second
// Ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){    return v > eps ? 1 : ( v < -eps ? -1
     : 0); }
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)
     ) <= 0; }
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
       p2) || btw(p4, p1, p2))
    return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&</pre>
    ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
ld a124 = cross(p2 - p1, p4 - p1);
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
     p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
(4))return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
       pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
```

dot(dp, dq));

4.2 Polar Angle Comparator [808e89]

4.3 Minkowski Sum [98abff]

```
void reorder_poly(vector<pdd>& pnts){
  int mn = 0;
  for(int i = 1; i < (int)pnts.size(); i++)</pre>
    if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y == pnts[mn</pre>
        ].Y && pnts[i].X < pnts[mn].X))
  rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
}
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
  reorder_poly(P);
  reorder_poly(Q);
  int psz = P.size();
  int qsz = Q.size();
  P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
  vector<pdd> ans;
  int i = 0, j = 0;
  while(i < psz || j < qsz){</pre>
    ans.pb(P[i] + Q[j]);
    int t = sgn(cross(P[i + 1] - P[i], Q[j + 1] - Q[j])
    if(t >= 0) i++;
    if(t <= 0) j++;
  }
  return ans;
}
```

4.4 Intersection of Circle and Convex Polygon

```
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
 if(a > r){
   S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
   S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
 else S = .5*sin(C)*a*b;
 return S;
double areaPolyCircle(const vector<pdd> poly,const pdd
    &0,const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly)]-0,r)*ori(0,
        poly[i],poly[(i+1)%SZ(poly)]);
 return fabs(S);
```

4.5 Intersection of Circles [f7a2fe]

```
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
```

4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
    int sign1 ){
  vector<Line> ret;
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
        v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sgn(p1.X - p2.X) == 0 and
        sgn(p1.Y - p2.Y) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  return ret;
}
```

4.7 Intersection of Line and Convex Polygon [157258]

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
  });
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 \mid | cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int 1, int r) {
    for (int t = 1; (1 + 1) % n != r; ) {
      int m = ((1 + r + (1 < r? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(t) ? 1 : r) = m;
    return (1 + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
  return pii(res.X, -1);
  if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
      +1
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned in the same order as the line hits the
      convex */
  return res;
} // convex cut: (r, l]
```

4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
    a, pdd b) {
    pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
    :
```

4.14

ConvexHull3D [156311]

struct convex_hull_3D {

```
National Taiwan University
  double s = cross(b - a, c - a), h2 = r * r - s * s /
      abs2(b - a);
  if (sgn(h2) < 0) return {};</pre>
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return \{p - h, p + h\};
4.9 Point in Circle [ecf954]
// return q's relation with circumcircle of tri(p[0],p
    [1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    int128 det = 0;
  for (int i = 0; i < 3; ++i)
    det += \underline{-int128(abs2(p[i]) - abs2(q)) * cross(p[(i
        + 1) % 3] - q, p[(i + 2) % 3] - q);
  return det > 0; // in: >0, on: =0, out: <0
4.10 Point in Convex [f86640]
bool PointInConvex(const vector<pll> &C, pll p, bool
    strict = true) {
  int a = 1, b = SZ(C) - 1, r = !strict;
  if (SZ(C) == 0) return false;
  if (SZ(C) < 3) return r && btw(C[0], C.back(), p);</pre>
  if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
  if (ori(C[0], C[a], p) >= r || ori(C[0], C[b], p) <=</pre>
      -r)
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (ori(C[0], C[c], p) > 0 ? b : a) = c;
  }
  return ori(C[a], C[b], p) < r;</pre>
}
4.11 Half Plane Intersection [dfb833]
// from 8BQube
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y - a.X, b.X - a.X), cross(a.Y - a
    .X, b.Y - a.X)); }
bool isin(Line 10, Line 11, Line 12) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(10, 12);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;</pre>
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0; // C^4
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
  sort(iter(arr), [&](Line a, Line b) -> int {
    if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
```

```
return cmp(a.Y - a.X, b.Y - b.X, 0);
 return ori(a.X, a.Y, b.Y) < 0;</pre>
deque<Line> dq(1, arr[0]);
for (auto p : arr) {
 if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) ==
       -1)
    continue;
 while (SZ(dq) >= 2 \&\& !isin(p, dq[SZ(dq) - 2], dq.
      back()))
    dq.pop_back();
 while (SZ(dq) >= 2 \&\& !isin(p, dq[0], dq[1]))
   dq.pop_front();
 dq.pb(p);
while (SZ(dq) >= 3 \&\& !isin(dq[0], dq[SZ(dq) - 2], dq
    .back()))
 dq.pop_back();
while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
 dq.pop_front();
return vector<Line>(iter(dq));
```

4.12 Minimum Enclosing Circle [5af6d5]

```
using ld = long double;
pair<pdd, ld> circumcenter(pdd a, pdd b, pdd c);
```

```
pair<pdd, ld> MinimumEnclosingCircle(vector<pdd> &pts){
  random_shuffle(iter(pts));
  pdd c = pts[0];
  1d r = 0;
  for(int i = 1; i < SZ(pts); i++){</pre>
    if(abs(pts[i] - c) <= r) continue;</pre>
    c = pts[i]; r = 0;
    for(int j = 0; j < i; j++){</pre>
      if(abs(pts[j] - c) <= r) continue;</pre>
      c = (pts[i] + pts[j]) / 2;
      r = abs(pts[i] - c);
      for(int k = 0; k < j; k++){
        if(abs(pts[k] - c) > r)
          tie(c, r) = circumcenter(pts[i], pts[j], pts[
    }
  return {c, r};
4.13 3D Point [badbbd]
// Copy from 8BQube
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x
      (_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
     pi1
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
    .y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
{ return (a + b + c + d) / 4; }
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point rotate_around(Point p, double angle, Point axis)
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
  return u * dot(u, p) * (1 - c) + p * c + cross(u, p)
}
```

```
struct Face {
                                                                triangulation such that no points will strictly
  int a, b, c;
                                                                inside circumcircle of any triangle. */
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
                                                            struct Edge {
}; // return the faces with pt indexes
                                                               int id; // oidx[id]
                                                              list<Edge>::iterator twin;
vector<Face> res:
vector<Point> P;
                                                              Edge(int _id = 0):id(_id) {}
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
                                                            struct Delaunay { // 0-base
  int n = SZ(P);
                                                              int n;
  if (n <= 2) return; // be careful about edge case</pre>
                                                              vector<int> oidx;
  // ensure first 4 points are not coplanar
                                                              vector<list<Edge>> head; // result udir. graph
  swap(P[1], *find_if(iter(P), [&](auto p) { return sgn
                                                              vector<pll> p;
      (abs2(P[0] - p)) != 0; }));
                                                              Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
  swap(P[2], *find_if(iter(P), [&](auto p) { return sgn
                                                                   head(n), p(n) {
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
                                                                 iota(iter(oidx), 0);
  swap(P[3],\ *find\_if(iter(P),\ [\&](auto\ p)\ \{\ return\ sgn
                                                                 for (int i = 0; i < n; ++i) head[i].clear();</pre>
      (volume(P[0], P[1], P[2], p)) != 0; }));
                                                                 sort(iter(oidx), [&](int a, int b)
  vector<vector<int>> flag(n, vector<int>(n));
                                                                     { return _p[a] < _p[b]; });
                                                                 for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>
  res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);
  for (int i = 3; i < n; ++i) {</pre>
                                                                 divide(0, n - 1);
    vector<Face> next;
    for (auto f : res) {
                                                              void addEdge(int u, int v) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
                                                                 head[u].push_front(Edge(v));
                                                                 head[v].push_front(Edge(u));
      if (d <= 0) next.pb(f);</pre>
                                                                 head[u].begin()->twin = head[v].begin();
      int ff = (d > 0) - (d < 0);
                                                                 head[v].begin()->twin = head[u].begin();
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
                                                               void divide(int 1, int r) {
                                                                 if (1 == r) return;
                                                                 if (1 + 1 == r) return addEdge(1, 1 + 1);
    for (auto f : res) {
                                                                 int mid = (1 + r) >> 1, nw[2] = \{1, r\};
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 && flag[y][x] <= 0)</pre>
                                                                 divide(l, mid), divide(mid + 1, r);
          next.emplace_back(x, y, i);
                                                                 auto gao = [&](int t)
                                                                   pll pt[2] = {p[nw[0]], p[nw[1]]};
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
                                                                   for (auto it : head[nw[t]]) {
                                                                     int v = ori(pt[1], pt[0], p[it.id]);
                                                                     if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
    res = next:
 }
                                                                         id]) < abs2(pt[1] - pt[0])))
                                                                       return nw[t] = it.id, true;
bool same(Face s, Face t) {
                                                                   }
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
                                                                   return false;
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
                                                                 while (gao(0) || gao(1));
                                                                 addEdge(nw[0], nw[1]); // add tangent
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
                                                                 while (true) {
       return 0;
                                                                   pll pt[2] = {p[nw[0]], p[nw[1]]};
                                                                   int ch = -1, sd = 0;
  return 1;
                                                                   for (int t = 0; t < 2; ++t)</pre>
int polygon_face_num() {
                                                                     for (auto it : head[nw[t]])
  int ans = 0;
                                                                       if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
  for (int i = 0; i < SZ(res); ++i)</pre>
                                                                            -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
    ans += none_of(res.begin(), res.begin() + i, [&](
                                                                           id])))
                                                                         ch = it.id, sd = t;
       Face g) { return same(res[i], g); });
  return ans;
                                                                   if (ch == -1) break; // upper common tangent
                                                                   for (auto it = head[nw[sd]].begin(); it != head[
double get_volume() {
                                                                       nw[sd]].end(); )
  double ans = 0;
                                                                     if (seg_strict_intersect(pt[sd], p[it->id], pt[
  for (auto f : res)
                                                                         sd ^ 1], p[ch]))
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
                                                                       head[it->id].erase(it->twin), head[nw[sd]].
        ]);
                                                                           erase(it++);
                                                                     else ++it;
  return fabs(ans / 6);
                                                                   nw[sd] = ch, addEdge(nw[0], nw[1]);
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
                                                              }
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
                                                            };
      .z) * (p3.y - p1.y);
                                                            4.16 Voronoi Diagram [e4f408]
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
  .x) * (p3.z - p1.z);

double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
                                                            // all coord. is even, you may want to call
                                                                 halfPlaneInter after then
      .y) * (p3.x - p1.x);
                                                            vector<vector<Line>> vec;
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
                                                             void build_voronoi_line(int n, vector<pll> &pts) {
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
                                                              Delaunay tool(n, pts); // Delaunay
       * a + b * b + c * c);
                                                              vec.clear(), vec.resize(n);
                                                              for (int i = 0; i < n; ++i)</pre>
                                                                 for (auto e : tool.head[i]) {
};
// n^2 delaunay: facets with negative z normal of
                                                                   int u = tool.oidx[i], v = tool.oidx[e.id];
// convexhull of (x, y, x^2 + y^2), use a pseudo-point // (0, 0, \inf) to avoid degenerate case
                                                                   pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v])
                                                                        - pts[u]);
                                                                   vec[u].pb(Line(m, m + d));
4.15 Delaunay Triangulation [6a9916]
/* Delaunay Triangulation:
                                                            }
   Given a sets of points on 2D plane, find a
```

4.17 Polygon Union [9fbf66]

```
// from 8BQube
ld rat(pll a, pll b) {
 return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
 // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
 ld res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {</pre>
          pll C = q[b], D = q[(b + 1) \% SZ(q)];
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            1d sa = cross(D - C, A - C), sb = cross(D -
                C, B - C);
            segs.pb(sa / (sa - sb), sgn(sc - sd));
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D
               - C)) > 0) {
            segs.pb(rat(C - A, B - A), 1);
            segs.pb(rat(D - A, B - A), -1);
          }
        }
      sort(iter(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[i].Y;
      res += cross(A, B) * sum;
    }
 return res / 2;
```

4.18 Tangent Point to Convex Hull [523bc1]

```
// from 8BQube
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
      { return ori(p, C[x], C[y]) == s; });
  };
  return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

4.19 Heart [082d19]

```
pdd circenter(pdd p0, pdd p1, pdd p2) { // radius = abs
    (center)
  p1 = p1 - p0, p2 = p2 - p0;
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
  pdd center;
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
     y1 - y2)) / m;
  center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
      y2 * y2) / m;
  return center + p0;
pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
     / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
 return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
     p3) * 2; }
```

4.20 Rotating Sweep Line [f5f689]

```
struct Event {
  pll d; int u, v;
  bool operator<(const Event &b) const {</pre>
    int ret = cmp(d, b.d, false);
    return ret == -1 ? false : ret; } // no tie-break
void rotatingSweepLine(const vector<pll> &p) {
  const int n = SZ(p);
  vector<Event> e; e.reserve(n * (n - 1));
  for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++) // pos[i] < pos[j] when</pre>
          the event occurs
      if (i != j) e.pb(p[j] - p[i], i, j);
  sort(iter(e));
  vector<int> ord(n), pos(n);
  iota(iter(ord), 0);
  sort(iter(ord), [&](int i, int j) { // initial order
      return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
           X < p[j].X; \});
  for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
  // initialize
  for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
    // do something
    vector<pii> tmp;
    for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
      tmp.pb(pii(e[j].u, e[j].v));
    sort(iter(tmp), [&](pii x, pii y){
        return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
            ], pos[y.ss]); });
    for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
      tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
        make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
             pos[x]);
}
```

4.21 Vector In Poly [c6d0fa]

```
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
    a-c
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
   return ori(a, b, p) >= strict && ori(a, p, c) >=
        strict;
}

// whether vector{cur, p} in counter-clockwise order
        prv, cur, nxt
bool inside(pll prv, pll cur, pll nxt, pll p, int
        strict) {
   if (ori(cur, nxt, prv) >= 0)
        return btwangle(cur, nxt, prv, p, strict);
   return !btwangle(cur, prv, nxt, p, !strict);
}
```

4.22 Convex Hull DP [92fd4b]

```
sort(iter(pts), [&](pll x, pll y) {
    return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;</pre>
    });
auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
// DP for convex hull vertices (no points on edges)
auto solve = [\&](int bottom) { // <math>O(n^3)
  pll 0 = pts[bottom];
  vector<pii> trans;
  for (int j = bottom + 1; j < n; j++)</pre>
    for (int k = bottom + 1; k < n; k++) {</pre>
      if (ori(0, pts[j], pts[k]) <= 0) continue;</pre>
      // check whether j->k is legal
      trans.pb(pii(j, k));
  sort(iter(trans), [&](pii x, pii y) -> bool{
      int tmp = cmp(getvec(x), getvec(y), false);
      if (tmp != -1) return tmp;
      pll v = getvec(x);
      return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
      });
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {
    // check whether bottom -> j is legal
    // init trans -> j
```

if(in[v] != -1){

continue;

}

if(in[v] < in[now]) up[now].pb(e);</pre>

low[now] = min(low[now], in[v]);

```
for (auto [i, j] : trans) {
                                                                     dfs(dfs, v, now, e);
   // normal trans i -> j
                                                                    low[now] = min(low[now], low[v]);
                                                                  if((now != par && low[now] >= in[par]) || (now ==
  for (int j = bottom + 1; j < n; j++) {</pre>
   // check whether j -> bottom is legal
                                                                        par && SZ(g[now]) == 0)){
                                                                     make_bcc();
    // end trans j ->
                                                                     for(int v = stk.back();; v = stk.back()){
                                                                       stk.pop_back(), add_v(v);
for(int i = 0; i < n; i++) solve(i);</pre>
                                                                       for(int e : up[v]) add_e(e);
                                                                       if(v == now) break;
4.23 Calculate Points in Triangle [bf746f]
                                                                     if(now != par) add_v(par);
// all points are distinct
                                                                  }
// cnt[i][j] = # of point k s.t. strictly above ij, and
     i < k < j
                                                                 for(int i = 0; i < n; i++)</pre>
// cnt2[i][j] = # of points k s.t. strictly in ij
                                                                  if(in[i] == -1) _dfs(_dfs, i, i, -1);
// preprocess space: O(n^2), time: O(n^3), query time:
                                                            };
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
                                                            5.2 SCC [2c9a01]
  for (int j = 0; j < n; j++){</pre>
    if (pts[i] >= pts[j]) continue;
                                                            struct SCC{ // 0-based, output reversed topo order
    for (int k = 0; k < n; k++) {
                                                              int n, cnt = 0;
      if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
                                                              vector<vector<int>> g;
        int tmp = ori(pts[i], pts[j], pts[k]);
                                                              vector<int> sccid;
        if (tmp > 0) cnt[i][j]++; // only for i < j</pre>
                                                              explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
                                                              void add_edge(int u, int v){
                                                                g[u].pb(v);
   }
                                                              void build(){
auto calc_tri = [&](array<int, 3> arr) { // strictly
                                                                vector<int> in(n, -1), low(n), stk;
    inside
                                                                vector<bool> instk(n);
  sort(iter(arr), [\&](int x, int y){ return pts[x] < }
                                                                int ts = 0;
      pts[y]; });
                                                                 auto dfs1 = [&](auto dfs, int now) -> void{
  auto [x, y, z] = arr;
                                                                  stk.pb(now); instk[now] = true;
  int tmp = ori(pts[x], pts[y], pts[z]);
                                                                   in[now] = low[now] = ts++;
  if (tmp == 0) return 0;
                                                                   for(int i : g[now]){
  else if (tmp < 0)</pre>
                                                                    if(in[i] == -1)
    return cnt[x][z] - cnt[x][y] - cnt[y][z] - cnt2[x][
                                                                       dfs(dfs, i), low[now] = min(low[now], low[i])
        y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
                                                                     else if(instk[i] && in[i] < in[now])</pre>
      x][z];
                                                                       low[now] = min(low[now], in[i]);
};
                                                                   if(low[now] == in[now]){
5
     Graph
                                                                     for(; stk.back() != now; stk.pop_back())
                                                                       sccid[stk.back()] = cnt, instk[stk.back()] =
5.1 BCC [d04ebe]
                                                                     sccid[now] = cnt++, instk[now] = false, stk.
struct BCC{ // O-based, allow multi edges but not allow
                                                                         pop_back();
     Loops
                                                                  }
  int n, m, cnt = 0;
                                                                };
  // n:|V|, m:|E|, cnt:#bcc
                                                                 for(int i = 0; i < n; i++)</pre>
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
                                                                  if(in[i] == -1) dfs1(dfs1, i);
  vector<vector<int>> bcc_v, bcc_e;
  vector<vector<pii>>> g; // original graph
                                                            };
  vector<pii> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
                                                            5.3 2-SAT [0686a5]
    n(_n), m(SZ(_edges)), g(_n), edges(_edges){
      for(int i = 0; i < m; i++){</pre>
                                                            struct SAT { // 0-based
        auto [u, v] = edges[i];
                                                              int n;
        g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
                                                              vector<bool> istrue;
      }
                                                              SCC scc;
                                                              SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
  void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
                                                              int neg(int a) {
  // modify these if you need more information
                                                                return a >= n ? a - n : a + n;
  void add_v(int v){ bcc_v.back().pb(v); }
  void add_e(int e){ bcc_e.back().pb(e); }
                                                              void add_clause(int a, int b) {
  void build(){
                                                                scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
    vector\langle int \rangle in(n, -1), low(n, -1), stk;
    vector<vector<int>> up(n);
                                                              bool solve() {
    int ts = 0;
                                                                scc.build();
    auto _dfs = [&](auto dfs, int now, int par, int pe)
                                                                 for (int i = 0; i < n; ++i) {</pre>
          -> void{
                                                                  if (scc.sccid[i] == scc.sccid[i + n]) return
      if(pe != -1) up[now].pb(pe);
                                                                       false;
      in[now] = low[now] = ts++;
                                                                   istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>
      stk.pb(now);
                                                                  istrue[i + n] = !istrue[i];
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
                                                                return true;
```

}

Dominator Tree [2da9bb]

};

```
struct Dominator {
 int n:
  vector<vector<int>> g, r, rdom; int tk;
  vector<int> dfn, rev, fa, sdom, dom, val, rp;
  Dominator(\textbf{int } \_n) \ : \ n(\_n), \ g(n), \ r(n), \ rdom(n), \ tk(0)
    dfn = rev = fa = sdom = dom =
      val = rp = vector<int>(n, -1); }
  void add_edge(int x, int y) { g[x].push_back(y); }
  void dfs(int x) {
    rev[dfn[x] = tk] = x;
    fa[tk] = sdom[tk] = val[tk] = tk; tk++;
    for (int u : g[x]) {
      if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
      r[dfn[u]].push_back(dfn[x]);
    }
  void merge(int x, int y) { fa[x] = y; }
  int find(int x, int c = 0) {
    if (fa[x] == x) return c ? -1 : x;
    if (int p = find(fa[x], 1); p != -1) {
      if (sdom[val[x]] > sdom[val[fa[x]]])
        val[x] = val[fa[x]];
      fa[x] = p;
      return c ? p : val[x];
    } else return c ? fa[x] : val[x];
  }
  vector<int> build(int s) {
    // return the father of each node in dominator tree dfs(s); // p[i] = -2 if i is unreachable, par[s] =
    for (int i = tk - 1; i >= 0; --i) {
      for (int u : r[i])
        sdom[i] = min(sdom[i], sdom[find(u)]);
      if (i) rdom[sdom[i]].push_back(i);
      for (int u : rdom[i]) {
        int p = find(u);
        dom[u] = (sdom[p] == i ? i : p);
      if (i) merge(i, rp[i]);
    }
    vector < int > p(n, -2); p[s] = -1;
    for (int i = 1; i < tk; ++i)</pre>
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
    for (int i = 1; i < tk; ++i)</pre>
      p[rev[i]] = rev[dom[i]];
    return p;
 }
5.5 Virtual Tree [6abeb5]
```

```
// copy from 8BQube
vector<int> vG[N];
int top, st[N];
int vrt = -1;
void insert(int u) {
  if (top == -1) return st[++top] = vrt = u, void();
  int p = LCA(st[top], u);
   if(dep[vrt] > dep[p]) vrt = p;
  if (p == st[top]) return st[++top] = u, void();
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
   vG[st[top - 1]].pb(st[top]), --top;
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
  st[++top] = u;
void reset(int u) {
  for (int i : vG[u]) reset(i);
  vG[u].clear();
void solve(vector<int> &v) {
  top = -1;
  sort(iter(v),
      [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
  for (int i : v) insert(i);
 while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
  // do something
  reset(vrt);
```

5.6 Fast DMST [7b274d]

```
struct E { int s, t; ll w; }; // 0-base
struct PQ {
  struct P {
    11 v; int i;
    bool operator>(const P &b) const { return v > b.v;
  priority_queue<P, vector<P>, greater<>> pq; 11 tag;
      // min heap
  void push(P p) { p.v -= tag; pq.emplace(p); }
  P top() { P p = pq.top(); p.v += tag; return p; }
  void join(PQ &b) {
    if (pq.size() < b.pq.size())</pre>
      swap(pq, b.pq), swap(tag, b.tag);
    while (!b.pq.empty()) push(b.top()), b.pq.pop();
}; // O(E log^2 V), use leftist tree for O(E log V)
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<PQ> h(n * 2);
  for (int i = 0; i < int(e.size()); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  vector<int> a(n * 2); iota(iter(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  auto o = [\&](auto Y, int x) \rightarrow int {
    return x==a[x] ? x : a[x] = Y(Y, a[x]); };
  auto S = [&](int i) { return o(o, e[i].s); };
  int pc = v[root] = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
    for (int p = i; v[p]<0 \mid \mid v[p]==i; p = S(r[p])) {
      if (v[p] == i)
        for (int q = pc++; p != q; p = S(r[p])) {
          h[p].tag -= h[p].top().v; h[q].join(h[p]);
          pa[p] = a[p] = q;
      while (S(h[p].top().i) == p) h[p].pq.pop();
      v[p] = i; r[p] = h[p].top().i;
  vector<int> ans;
  for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
    for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
         1)
      v[f] = n;
    ans.push_back(r[i]);
  return ans; // default minimize, returns edgeid array
}
```

```
5.7 Vizing [58a6ca]
// find D+1 edge coloring of a graph with max deg D, O(
              nm)
struct Vizing { // returns maxdeg+1 edge coloring in
              adjacent matrix G
       int n; // 1-based for vertices and colors, simple
                     graph
       vector<vector<int>> C, G;
       vector<int> X, vst;
       Vizing(int _n): n(_n),
       C(n + 1, vector < int > (n + 2)), G(n + 1, vector < int > (n + 1, 
                    + 1)),
       X(n + 1, 1), vst(n + 1) {}
       void solve(vector<pii> &E) {
              auto update = [&](int u)
              { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
              auto color = [&](int u, int v, int c) {
                    int p = G[u][v];
                    G[u][v] = G[v][u] = c;
                    C[u][c] = v, C[v][c] = u;
                    C[u][p] = C[v][p] = 0;
                    if (p) X[u] = X[v] = p;
                     else update(u), update(v);
                    return p;
              };
              auto flip = [&](int u, int c1, int c2) {
                    int p = C[u][c1];
                     swap(C[u][c1], C[u][c2]);
                    if (p) G[u][p] = G[p][u] = c2;
                    if (!C[u][c1]) X[u] = c1;
                    if (!C[u][c2]) X[u] = c2;
                     return p;
              };
```

```
for (int t = 0; t < SZ(E); ++t) {</pre>
      int u = E[t].ff, v0 = E[t].ss, v = v0, c0 = X[u],
           c = c0, d;
      vector<pii> L;
      fill(iter(vst), 0);
      while (!G[u][v0]) {
        L.emplace_back(v, d = X[v]);
        if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
             --a) c = color(u, L[a].ff, c);
        else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
             0; --a) color(u, L[a].ff, L[a].ss);
        else if (vst[d]) break;
        else vst[d] = 1, v = C[u][d];
      if (!G[u][v0]) {
        for (; v; v = flip(v, c, d), swap(c, d));
        if (int a; C[u][c0]) {
          for (a = SZ(L) - 2; a >= 0 \&\& L[a].ss != c;
              --a);
          for (; a >= 0; --a) color(u, L[a].ff, L[a].ss
              );
        else --t:
     }
 }
};
     Maximum Clique [1ad4b2]
  bitset<N> G[N], cs[N];
```

5.8

```
struct MaxClique { // fast when N <= 100</pre>
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
    G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(iter(r), [\&](int x, int y) \{ return d[x] > d \}
          [y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < 1ft) r[tp++] = p;
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
          ]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<int> &r, vector<int> &c, int 1,
      bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(iter(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
};
```

Number of Maximal Clique [11fa26] 5.9

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
      int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
        if (g[v][some[d][j]])
          some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
        if (g[v][none[d][j]])
          none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
};
```

Minimum Mean Cycle [3e5d2b] 5.10

```
// from 8BOube
11 road[N][N]; // input here
struct MinimumMeanCycle {
  11 dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
for (int i = 2; i <= L; ++i)
       for (int k = 0; k < n; ++k)
         for (int j = 0; j < n; ++j)</pre>
           dp[i][j]
             min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
       if (dp[L][i] >= INF) continue;
       11 ta = 0, tb = 1;
       for (int j = 1; j < n; ++j)</pre>
         if (dp[j][i] < INF &&
    ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)</pre>
           ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
       if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
       11 g = \_gcd(a, b);
       return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
};
```

Minimum Steiner Tree [21acea] 5.11

```
// from 8BQube
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];</pre>
```

```
int vcost[N]; // the cost of vertexs
  void init(int _n) {
   n = _n;
    for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
   }
  }
  void add_edge(int ui, int vi, int wi) {
   dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)</pre>
        for (int j = 0; j < n; ++j)</pre>
          dst[i][j] =
            min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
        for (int i = 0; i < n; ++i)</pre>
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
             submsk = (submsk - 1) \& msk)
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
              vcost[i]);
      for (int i = 0; i < n; ++i) {
        tdst[i] = INF;
        for (int j = 0; j < n; ++j)</pre>
          tdst[i] =
            min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    }
    int ans = INF;
    for (int i = 0; i < n; ++i)</pre>
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
 }
};
5.12 Count Cycles [c7e8f2]
```

```
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
  for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
 for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) c4 += vis[z]++;
  for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M))
```

Math

ll ifloor(ll a, ll b){

6.1 Extended Euclidean Algorithm [c51ae9]

```
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
void extgcd(ll x, ll y, ll &g, ll &a, ll &b) {
 if (y == 0) g = x, a = 1, b = 0;
  else extgcd(y, x % y, g, b, a), b -= (x / y) * a;
6.2 Floor & Ceil [134881]
```

return a / b - (a % b && (a < 0) ^ (b < 0));

```
15
11 iceil(11 a, 11 b){
 return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre [4e4b23]
// the Jacobi symbol is a generalization of the
    Legendre symbol,
// such that the bottom doesn't need to be prime.
// (n|p) -> same as Legendre
// (n|ab) = (n|a)(n|b)
// work with long long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
    a >>= r;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  }
  return s;
}
// 0: a == 0
// -1: a isn't a quad res of p
// else: return X with X^2 % p == a
// doesn't work with long long
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
  int b, d;
  for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
    if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
           p)) % p;
      g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
      g0 = tmp;
    tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
    )) % p;
f1 = (2LL * f0 * f1) % p;
    f0 = tmp;
  }
  return g0;
}
6.4 Simplex [aa7741]
 // maximize c^T x
// subject to Ax <= b, x >= 0
// and stores the solution;
typedef long double T; // long double, Rational, double
     + mod<P>..
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-9, inf = 1/.0;
#define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
    s])) s=j
#define rep(i, l, n) for(int i = l; i < n; i++)
struct LPSolver {
  int m, n;
  vector<int> N, B;
  vvd D;
  LPSolver(const vvd& A, const vd& b, const vd& c) :
    m(SZ(b)), n(SZ(c)), N(n+1), B(m), D(m+2, vd(n+2)) {
      rep(i,0,m) \ rep(j,0,n) \ D[i][j] = A[i][j];
      rep(i,0,m) { B[i] = n+i; D[i][n] = -1; D[i][n+1]
          = b[i];}
      rep(j,0,n) \{ N[j] = j; D[m][j] = -c[j]; \}
```

N[n] = -1; D[m+1][n] = 1;

```
if (s == y) return 100;
  void pivot(int r, int s) {
                                                                     int p = 100 + DiscreteLog(s, x, y, m);
     T *a = D[r].data(), inv = 1 / a[s];
                                                                     if (fpow(x, p, m) != y) return -1;
     rep(i,0,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
                                                                     return p; //returns: x^p = y \pmod{m}
       T *b = D[i].data(), inv2 = b[s] * inv;
       rep(j,0,n+2) b[j] -= a[j] * inv2;
                                                                   6.7 Miller Rabin & Pollard Rho [d3ecd2]
       b[s] = a[s] * inv2;
                                                                   // n < 4,759,123,141
                                                                                                 3 : 2, 7, 61
     rep(j,0,n+2) if (j != s) D[r][j] *= inv;
                                                                   // n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : primes <= 13
    rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
                                                                   // n < 2^64
     swap(B[r], N[s]);
                                                                   // 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                   11 mul(ll a, ll b, ll n){
                                                                     return (__int128)a * b % n;
  bool simplex(int phase) {
    int x = m + phase - 1;
                                                                   bool Miller_Rabin(ll a, ll n) {
   if ((a = a % n) == 0) return 1;
     for (;;) {
       int s = -1;
                                                                     if (n % 2 == 0) return n == 2;
       rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
                                                                     ll tmp = (n - 1) / ((n - 1) & (1 - n));
ll t = __lg(((n - 1) & (1 - n))), x = 1;
       if (D[x][s] >= -eps) return true;
       int r = -1;
                                                                     for (; tmp; tmp >>= 1, a = mul(a, a, n))
       rep(i,0,m) {
                                                                       if (tmp \& 1) x = mul(x, a, n);
         if (D[i][s] <= eps) continue;</pre>
                                                                     if (x == 1 || x == n - 1) return 1;
         if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])
                                                                     while (--t)
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
                                                                       if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                     return 0;
       if (r == -1) return false;
       pivot(r, s);
                                                                   bool prime(ll n){
     }
                                                                     vector<ll> tmp = \{2, 325, 9375, 28178, 450775,
                                                                          9780504, 1795265022};
                                                                     for(ll i : tmp)
  T solve(vd &x) {
                                                                       if(!Miller_Rabin(i, n)) return false;
     int r = 0;
                                                                     return true;
     rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
     if (D[r][n+1] < -eps) {</pre>
                                                                   map<11, int> cnt;
       pivot(r, n);
                                                                   void PollardRho(ll n) {
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                     if (n == 1) return;
           inf;
                                                                     if (prime(n)) return ++cnt[n], void();
       rep(i,0,m) if (B[i] == -1) {
                                                                     if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
                                                                          void();
         rep(j,1,n+1) ltj(D[i]);
                                                                     11 x = 2, y = 2, d = 1, p = 1;
         pivot(i, s);
                                                                   #define f(x, n, p) ((mul(x, x, n) + p) % n)
       }
                                                                     while (true) {
                                                                       if (d != n && d != 1) {
     bool ok = simplex(1); x = vd(n);
                                                                          PollardRho(n / d);
     rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
                                                                          PollardRho(d);
     return ok ? D[m][n+1] : inf;
                                                                          return;
  }
                                                                       if (d == n) ++p;
                                                                       x = f(x, n, p), y = f(f(y, n, p), n, p);
       Simplex Construction
                                                                       d = gcd(abs(x - y), n);
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for
all 1 \leq j \leq m and x_i \geq 0 for all 1 \leq i \leq n.
1. In case of minimization, let c_i' = -c_i
                                                                   6.8 XOR Basis [006505]
2. \sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j
3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j \to \mathsf{add} \subseteq \mathsf{and} \supseteq .
                                                                   const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                   struct Basis{
6.6 DiscreteLog [da27bf]
                                                                     int total = 0, rank = 0;
                                                                     vector<ll> b;
int DiscreteLog(int s, int x, int y, int m) {
                                                                     Basis(): b(digit) {}
  constexpr int kStep = 32000;
                                                                     bool add(ll v){ // Gauss Jordan Elimination
  unordered_map<int, int> p;
                                                                       total++:
  int b = 1;
                                                                        for(int i = digit - 1; i >= 0; i--){
  for (int i = 0; i < kStep; ++i) {</pre>
                                                                          if(!(1LL << i & v)) continue;</pre>
    p[y] = i;
y = 1LL * y * x % m;
                                                                          if(b[i] != 0){
                                                                            v ^= b[i];
     b = 1LL * b * x % m;
                                                                            continue;
  for (int i = 0; i < m + 10; i += kStep) {</pre>
                                                                          for(int j = 0; j < i; j++)</pre>
    s = 1LL * s * b % m;
                                                                            if(1LL << j & v) v ^= b[j];</pre>
     if (p.find(s) != p.end()) return i + kStep - p[s];
                                                                          for(int j = i + 1; j < digit; j++)</pre>
                                                                            if(1LL << i & b[j]) b[j] ^= v;</pre>
  return -1;
                                                                          b[i] = v;
                                                                          rank++;
int DiscreteLog(int x, int y, int m) {
                                                                          return true;
  if (m == 1) return 0;
  int s = 1;
                                                                       return false;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
                                                                     11 \text{ getmax}(11 \text{ x} = 0){
     s = 1LL * s * x % m;
```

for(ll i : b) $x = max(x, x ^ i);$

return x;

ll lead = mat[rk][i];

* inv(lead) % MOD;

for (int j = 0; j < N; j++) {</pre>

for (int k = 0; k < M; k++)

+ MOD) % MOD;

vector<11> par; //particular solution (Ax = b)

vector<vector<ll>>> homo; //homogenous (Ax = 0)

void solve(const vector<vector<11>>> &eq) {

if(piv.size() && piv.back() == M)

if (j == rk) continue;

struct LinearEquation { // 2702e2

//first M columns are matrix A

//last column of eq is vector b

vector<vector<ll>> rref;

int M = SZ(eq[0]) - 1;

auto piv = RREF(rref);

vector<bool> ispiv(M);

ispiv[piv[i]] = 1;

vector<11> h(M);

h[i] = 1;

return ok = 0, void();

for (int i = 0;i < rk;i++) {</pre> par[piv[i]] = rref[i][M];

for (int i = 0; i < M; i++) {</pre>

if (ispiv[i]) continue;

int rk = piv.size();

11 tmp = mat[j][i];

cols.pb(i);

rk++;

return cols;

rref = eq;

ok = 1:

par.resize(M);

```
National Taiwan University
  ll getmin(ll x = 0){
    for(ll i : b) x = min(x, x ^ i);
    return x;
  bool can(ll x){
    return getmin(x) == 0;
  11 kth(11 k){ // kth smallest, 0-indexed
    vector<ll> tmp;
    for(11 i : b) if(i) tmp.pb(i);
    11 \text{ ans} = 0;
    for(int i = 0; i < SZ(tmp); i++)</pre>
      if(1LL << i & k) ans ^= tmp[i];</pre>
    return ans;
 }
};
6.9
    Linear Equation [056191]
vector<int> RREF(vector<vector<11>> &mat) { // 9cd26b
  int N = SZ(mat), M = SZ(mat[0]);
  int rk = 0;
  vector<int> cols;
  for (int i = 0; i < M; i++) {</pre>
    int cnt = -1;
    for (int j = N - 1; j >= rk; j--)
      if(mat[j][i] != 0) cnt = j;
    if (cnt == -1) continue;
    swap(mat[rk], mat[cnt]);
```

for (int j = 0; j < M; j++) mat[rk][j] = mat[rk][j]</pre>

mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % MOD

for (int j = 0; j < rk; j++)</pre> h[piv[j]] = rref[j][i] ? MOD - rref[j][i] : 0; homo.pb(h); } } };

6.10 Chinese Remainder Theorem [6ef4a3]

```
pll solve crt(ll x1, ll m1, ll x2, ll m2){
 ll g = gcd(m1, m2);
  if ((x2 - x1) % g) return {0, 0}; // no sol
  m1 /= g; m2 /= g;
```

```
11 _, p, q;
   extgcd(m1, m2, _, p, q); // p <= C
ll lcm = m1 * m2 * g;
   ll res = ((__int128)p * (x2 - x1) % lcm * m1 % lcm +
          x1) % lcm;
   // be careful with overflow, C^3
   return {(res + lcm) % lcm, lcm}; // (x, m)
6.11 Sqrt Decomposition [8d7bc0]
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){</pre>
   int x = ifloor(n, 1);
   r = ifloor(n, x);
// for all i in [l, r], ceil(n / i) = x
for(int 1, r = n; r >= 1; r = 1 - 1){
   int x = iceil(n, r);
   l = iceil(n, x);
6.12 Floor Sum
• m = |\frac{an+b}{2}|
• Time complexity: O(\log n)
      f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
                           \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), \end{cases}
                                                                       a \ge c \lor b \ge c
                                                                        n < 0 \lor a = 0
                           nm - f(c, c - b - 1, a, m - 1), otherwise
g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor
                     \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                     +g(a \bmod c, b \bmod c, c, n),
                                                                               a \geq c \lor b \geq c
                                                                               n < 0 \lor a = 0
                     \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                    (-h(c,c-b-1,a,m-1)),
                                                                               otherwise
h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2
                     \left(\lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1)\right)
                     +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                     +h(a\ \mathsf{mod}\ c,b\ \mathsf{mod}\ c,c,n)
                    \int +2\lfloor rac{a}{c} 
floor \cdot g(a mod c, b mod c, c, n)
                     +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                               a \ge c \lor b \ge c
                     0,
                                                                               n < 0 \lor a = 0
                     nm(m+1) - 2g(c, c-b-1, a, m-1)
                    -2f(c,c-b-1,a,m-1)-f(a,b,c,n), otherwise
```

Polynomial 7

7.1 FWHT [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)] or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)</pre>
    for (int i = 0; i < n; i += L)
       for (int j = i; j < i + (L >> 1); ++j)
         a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
     1;
void subset_convolution(int *a, int *b, int *c, int L)
  // c_k = \sum_{i=0}^{n} \{i \mid j = k, i \& j = 0\} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
```

```
for (int i = 0; i <= L; ++i)</pre>
    fwt(f[i], n, 1), fwt(g[i], n, 1);
  for (int i = 0; i <= L; ++i)
    for (int j = 0; j <= i; ++j)</pre>
      for (int x = 0; x < n; ++x)
         h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i) fwt(h[i], n, -1);
for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
7.2 FFT [13ec2f]
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double
using val_t = complex<double>;
                                                                 };
template<int MAXN>
struct FFT {
  const double PI = acos(-1);
  val_t w[MAXN];
  FFT() {
    for (int i = 0; i < MAXN; ++i) {
  double arg = 2 * PI * i / MAXN;</pre>
      w[i] = val_t(cos(arg), sin(arg));
    }
  void bitrev(vector<val_t> &a, int n) //same as NTT
void trans(vector<val_t> &a, int n, bool inv = false)
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
         for (int j = i, x = 0; j < i + d1; ++j, x += dx
           val_t + mp = a[j + dl] * (inv ? conj(w[x]) : w
                [x]);
           a[j + d1] = a[j] - tmp;
           a[j] += tmp;
        }
      }
    if (inv) {
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
  //multiplying two polynomials A * B:
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
  //A[i] *= B[i], fft.trans(A, siz, 1);
7.3 NTT [bf683f]
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
// only works when sz(A) + sz(B) - 1 <= MAXN
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
struct NTT {
  11 w[MAXN];
  11 mpow(ll a, ll n);
  11 minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
          % P;
  void bitrev(vector<ll> &a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
    }
  }
  void operator()(vector<ll> &a, int n, bool inv =
       false) { //0 <= a[i] < P
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {
      int dx = MAXN / L, dl = L >> 1;
      for (int i = 0; i < n; i += L) {</pre>
         for (int j = i, x = 0; j < i + dl; ++j, x += dx
             ) {
```

```
11 \text{ tmp} = a[j + d1] * w[x] % P;
        if ((a[j + dl] = a[j] - tmp) < 0) a[j + dl]
             += P:
        if ((a[j] += tmp) >= P) a[j] -= P;
    }
  if (inv) {
    reverse(a.begin()+1, a.begin()+n);
    11 invn = minv(n);
    for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
}
   Polynomial Operation [77a8a8]
```

```
7.4
// == b4233a ==
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
    ++i)
#define neg(x) (x ? P - x : 0)
#define V (*this)
template <int MAXN, 11 P, 11 RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
  using vector<ll>::vector;
  static inline NTT<MAXN, P, RT> ntt;
  int n() const { return (int)size(); }
                                           // n() >= 1
  Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
       .data(), min(p.n(), m), data()); }
  Poly &irev() { return reverse(data(), data() + n()),
  Poly &isz(int m) { return resize(m), V; }
  static ll minv(ll x) { return ntt.minv(x); }
  == fb1867 ==
  Poly &iadd(const Poly &rhs) { // db5668
    fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
    return V; // need n() == rhs.n()
  Poly &imul(11 k) { // a8df26
    fi(0, n()) V[i] = V[i] * k % P;
  Poly Mul(const Poly &rhs) const { // 46caf3
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    assert(m <= MAXN);</pre>
    Poly X(V, m), Y(rhs, m);
    ntt(X, m), ntt(Y, m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X, m, true);
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // 796a37
    if (n() == 1) return {minv(V[0])};
    int m = 1; // need V[0] != 0, 2*sz<=MAXN</pre>
    while (m < n() * 2) m <<= 1;
    assert(m <= MAXN);</pre>
    Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);
    Poly Y(V, m);
    ntt(Xi, m), ntt(Y, m);
    fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
      if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    ntt(Xi, m, true);
    return Xi.isz(n());
  Poly &shift_inplace(const ll &c) { // 0c04f6
    int n = V.n(); // 2 * sz <= MAXN</pre>
    vector<ll> fc(n), ifc(n);
    fc[0] = ifc[0] = 1;
    for (int i = 1; i < n; i++) {</pre>
      fc[i] = fc[i - 1] * i % P;
      ifc[i] = minv(fc[i]);
    for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
    Poly g(n);
    11 cp = 1;
    for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
         cp = cp * c % P;
    V = V.irev().Mul(g).isz(n).irev();
```

```
for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
                                                                      Mul(up[i * 2 + 1]);
                                                                 return up;
    return V;
  == 7b2835 ==
  Poly shift(const 11 &c) const { return Poly(V).
      shift_inplace(c); }
                                                                 auto up = _tree1(x);
  Poly \_Sqrt() const { // Jacobi(V[0], P) = 1
                                                                 return _eval(x, up);
    if (n() == 1) return {QuadraticResidue(V[0], P)};
    Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
                                                                   vector<11> &y) { // d7bae4
    return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
  }
// == b46641 ==
  Poly Sqrt() const { // 1aa942
    Poly a; // 2 * sz <= MAXN
                                                                 fi(0, m) down[m + i] = {z[i]};
    bool has = 0;
    for (int i = 0; i < n(); i++) {</pre>
      if (V[i]) has = 1;
      if (has) a.push_back(V[i]);
                                                                 return down[1];
    if (!has) return V;
                                                             // == c066ab ==
    if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {
      return Poly();
    a = a.isz((n() + a.n()) / 2)._Sqrt();
    int sz = a.n();
                                                                 if (n() == 1) return {1};
    a.isz(n());
    rotate(a.begin(), a.begin() + sz, a.end());
                                                                 Poly Y = X.Ln();
                                                                 Y[0] = P - 1;
  pair<Poly, Poly> DivMod(const Poly &rhs) const { // 5
                                                                 return X.Mul(Y).isz(n());
    if (n() < rhs.n()) return {{0}, V};</pre>
                                                             // == 3f1d86 ==
    const int m = n() - rhs.n() + 1;
    Poly X(rhs); // (rhs.)back() != 0
    X.irev().isz(m);
                                                                 int nz = 0:
    Poly Y(V);
    Y.irev().isz(m);
    Poly Q = Y.Mul(X.Inv()).isz(m).irev();
    X = rhs.Mul(Q), Y = V;
    fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
return {0, Y.isz(max(1, rhs.n() - 1))};
                                                                      n()).irev();
  == 76b1af ==
  Poly Dx() const {
    Poly ret(n() - 1);
    fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
    return ret.isz(max(1, ret.n()));
                                                                 if (!n()) {
                                                                   return Poly(m + 1, 0);
  Poly Sx() const {
    Poly ret(n() + 1);
                                                                  if (V[0] != 0) {
                                                                   1\dot{1}\dot{c} = V[0];
    fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
    return ret;
                                                                   V[0] = 0;
  Poly _tmul(int nn, const Poly &rhs) const {
    Poly Y = Mul(rhs).isz(n() + nn - 1);
                                                                       factorial
    return Poly(Y.data() + n() - 1, Y.data() + Y.n());
                                                                    Poly B(m + 1);
                                                                   11 pow = 1;
// == 3afa3f ==
                                                                   * c % P; // inv. of fac
A = A.Mul(B).isz(m + 1);
  vector<ll> _eval(const vector<ll> &x, const vector<</pre>
      Poly> &up) const { // fb6553
    const int m = (int)x.size();
    if (!m) return {};
                                                                   return A;
    vector<Poly> down(m * 2);
    // down[1] = DivMod(up[1]).second;
    // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
                                                                 int n = 1;
                                                                 while (n < V.n()) n *= 2;</pre>
    down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
                                                                 isz(n), wt.isz(n).irev();
        _tmul(m, V);
                                                                 int k = 1;
    fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
        1, down[i / 2]);
                                                                 q.imul(P - 1);
    vector<11> y(m);
                                                                 while (n > 1) {
    fi(0, m) y[i] = down[m + i][0];
                                                                   Poly r(2 * n * k);
    return y;
  static vector<Poly> _tree1(const vector<ll> &x) { //
                                                                       q[i]));
      f5c433
    const int m = (int)x.size();
                                                                   fi(0, 2 * n * k) {
    vector<Poly> up(m * 2);
                                                                      pq[2 * n * k + i] += p[i];
    fi(0, m) up[m + i] = {neg(x[i]), 1};
```

```
for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
vector<ll> Eval(const vector<ll> &x) const { // 1e5,
static Poly Interpolate(const vector<11> &x, const
  const int m = (int)x.size(); // 1e5, 1.4s
  vector<Poly> up = _{tree1(x), down(m * 2);}
  vector<ll> z = up[1].Dx()._eval(x, up);
  fi(0, m) z[i] = y[i] * minv(z[i]) % P;
  for (int i = m - 1; i > 0; --i)
  down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
        down[i * 2 + 1].Mul(up[i * 2]));
Poly Ln() const \{ // V[\theta] == 1, 2*sz <= MAXN \}
  return Dx().Mul(Inv()).Sx().isz(n());
Poly Exp() const { //V[0] == 0,2*sz <= MAXN
  Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
  fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
//M := P(P - 1). If k >= M, k := k % M + M.
Poly Pow(11 k) const { // 2*sz<=MAXN // d08261
  while (nz < n() && !V[nz]) ++nz;</pre>
  if (nz * min(k, (ll)n()) >= n()) return Poly(n());
  if (!k) return Poly(Poly{1}, n());
  Poly X(data() + nz, data() + nz + n() - nz * k);
  const 11 c = ntt.mpow(X[0], k % (P - 1));
  return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
// sum_j w_j [x^j] f(x^i) for i \in [0, m]
Poly power_projection(Poly wt, int m) { // 277119
  assert(n() == wt.n()); // 4*sz <= MAXN!
    Poly A = V.power_projection(wt, m);
    fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
    fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
    fi(0, m + 1) A[i] = A[i] * fac[i] % P;
  Poly p(wt, 2 * n), q(V, 2 * n);
    fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
    Poly pq = p.Mul(r).isz(4 * n * k);
    Poly qq = q.Mul(r).isz(4 * n * k);
      qq[2 * n * k + i] += q[i] + r[i];
```

```
pq[2 * n * k + i] %= P;
      qq[2 * n * k + i] %= P;
    fill(p.begin(), p.end(), 0);
    fill(q.begin(), q.end(), 0);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
  p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];
  q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];</pre>
    n /= 2, k *= 2;
  Poly ans(k);
  fi(0, k) ans[i] = p[2 * i];
  return ans.irev().isz(m + 1);
Poly FPSinv() { // 2c54b4
  const int n = V.n() - 1;
  if (n == -1) return {};
  assert(V[0] == 0);
  if (n == 0) return V;
  assert(V[1] != 0);
  ll c = V[1], ic = minv(c);
  imul(ic);
  Poly wt(n + 1);
  wt[n] = 1;
  Poly A = V.power_projection(wt, n);
  Poly g(n);
  fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
  g = g.Pow(neg(minv(n)));
  g.insert(g.begin(), 0);
  fi(0, g.n()) g[i] = g[i] * pow % P, pow = pow * ic
      % P;
  return g;
Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
    ] = k; // 7b552c
  return Poly(*this).irev().Mul(rhs).isz(n()).irev();
Poly FPScomp(Poly g) { // solves V(g(x)) // 332bb2
  auto rec = [&](auto &rec, int n, int k, Poly Q) ->
    if (n == 1) {
      Poly p(2 * k);
      irev();
      fi(0, k) p[2 * i] = V[i];
      return p;
    Poly R(2 * n * k);
    fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
         Q[i]));
    Poly QQ = Q.Mul(R).isz(4 * n * k);
fi(0, 2 * n * k) {
      QQ[2 * n * k + i] += Q[i] + R[i];
      QQ[2 * n * k + i] %= P;
    Poly nxt_Q(2 * n * k);
    for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {

nxt_Q[n * j + i] = QQ[(2 * n) * j + (2 * i + 0)
    Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
Poly pq(4 * n * k);
for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
      pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +
      i];
pq[(2 * n) * j + (2 * i + 1)] %= P;
    Poly p(2 * n * k);
    fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i
         ]) % P;
    pq.pop_back();
    Poly x = pq.TMul(R);
    fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
    return p;
  int sz = 1;
  while(sz < n() || sz < g.n()) sz <<= 1;
  return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
        sz)).isz(sz).irev();
```

```
}
};
#undef fi
#undef V
#undef neg
using Poly_t = Poly<1 << 19, 998244353, 3>;
```

7.5 Generating Function

Ordinary Generating Function

• C(x)=A(rx): $c_n=r^na_n$ 的一般生成函數。 • C(x)=A(x)+B(x): $c_n=a_n+b_n$ 的一般生成函數。

• C(x) = A(x)B(x): $c_n = \sum\limits_{i=0}^n a_i b_{n-i}$ 的一般生成函數。

• $C(x)=A(x)^k$: $c_n=\sum_{i_1+i_2+\ldots+i_k=n}^{i=0}a_{i_1}a_{i_2}\ldots a_{i_k}$ 的一般生成函數。

• C(x) = xA(x)': $c_n = na_n$ 的一般生成函數。

• $C(x) = \frac{A(x)}{1-x}$: $c_n = \sum_{i=0}^n a_i$ 的一般生成函數。

• $C(x)=A(1)+x\frac{A(1)-A(x)}{1-x}$: $c_n=\sum\limits_{i=n}^{\infty}a_i$ 的一般生成函數。

常用展開式

- $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {a \choose n} x^n$, ${a \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

常見生函

• 卡特蘭數: $f(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$

Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數
- $\hat{C}(x)=\hat{A}(x)^k$: $\sum_{i_1+i_2+\dots+i_k=n}^{\sum_{k=0}^{n}\binom{n}{i_1,i_2,\dots,i_k}}a_ia_{i_2}\dots a_{i_k}$ 的指數生成函數
- $\hat{C}(x)=\exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 F(0),G(0)=0, $F'(0),G'(0)\neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$
$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

7.6 Bostan Mori [41c3bc]

```
const 11 mod = 998244353;
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
     k)
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = \{P, Q\};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  vector<11> Qn(sz);
  vector N(2, vector<11>(sz));
  while(k) {
    fill(iter(Qn), 0);
    for(int i = 0; i < d + 1; i++){
      Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    ntt(Qn, sz, false);
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){</pre>
      fill(iter(N[i]), 0);
      copy(iter(M[i]), N[i].begin());
      ntt(N[i], sz, false);
for(int j = 0; j < sz; j++)</pre>
```

```
N[i][j] = N[i][j] * Qn[j] % mod;
     ntt(N[i], sz, true);
     for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
       M[i][j >> 1] = N[i][j];
    k >>= 1;
 }
  return M[0][0] * ntt.minv(M[1][0]) % mod;
11 LinearRecursion(vector<11> a, vector<11> c, 11 k) {
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
     1)));
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
  auto A = a; A.resize(sz);
  auto C = c; C.resize(sz);
  ntt(A, sz, false), ntt(C, sz, false);
  for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
 ntt(A, sz, true);
 A.resize(d);
 return BostanMori(A, c, k);
    String
8.1 KMP Algorithm [c8b75f]
```

```
// fail[i] = max k < i s.t. s[0..k] = s[i-k..i]
vector<int> kmp_build_fail(const string &s){
  int n = SZ(s);
  vector<int> fail(n, -1);
  int cur = -1;
  for(int i = 1; i < n; i++){</pre>
    while(cur != -1 && s[cur + 1] != s[i])
      cur = fail[cur];
    if(s[cur + 1] == s[i])
      cur++;
    fail[i] = cur;
  }
  return fail;
void kmp_match(const string &s, const vector<int> &fail
    , const string &t){
  int cur = -1;
  int n = SZ(s), m = SZ(t);
  for(int i = 0; i < m; i++){</pre>
    while(cur != -1 \&\& (cur + 1 == n || s[cur + 1] != t
        [i]))
      cur = fail[cur];
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
      cur++:
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
}
```

8.2 Manacher Algorithm [caf0f4]

```
/* center i: radius z[i * 2 + 1] / 2
  center i, i + 1: radius z[i * 2 + 2] / 2
   both aba, abba have radius 2 */
vector<int> manacher(const string &tmp){ // 0-based
  string s = "%";
  int 1 = 0, r = 0;
  for(char c : tmp) s += c, s += '%';
  vector<int> z(SZ(s));
  for(int i = 0; i < SZ(s); i++){</pre>
   z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
    while(i - z[i] >= 0 \&\& i + z[i] < SZ(s)
           && s[i + z[i]] == s[i - z[i]])
      ++z[i];
   if(z[i] + i > r) r = z[i] + i, l = i;
 }
  return z;
```

8.3 Lyndon Factorization [7c612b]

8.4 Suffix Array [cd67ea]

```
struct SuffixArray {
  vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]
       and sa[i-1]
                                 // sa[0] = s.size()
                                 // character should be 1-
                                      based
  SuffixArray(string& s, int lim=256) { // or
       basic_string<int>
    int n = s.size() + 1, k = 0, a, b;
vector<int> x(n, 0), y(n), ws(max(n, lim));
    rank.assign(n, 0);
    for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
    sa = lcp = y, iota(sa.begin(), sa.end(), 0);
for (int j = 0, p = 0; p < n; j = max(1, j * 2),</pre>
         lim = p) {
       p = j, iota(y.begin(), y.end(), n - j);
       for (int i = 0; i < n; i++)
         if (sa[i] >= j) y[p++] = sa[i] - j;
       for (int &i : ws) i = 0;
       for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
       for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];</pre>
       for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
       swap(x, y), p = 1, x[sa[0]] = 0;
       for(int i = 1; i < n; i++){</pre>
         a = sa[i - 1], b = sa[i];
         x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ?
               p - 1 : p++;
      }
    for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
    for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
       for (k && k--, j = sa[rank[i] - 1];
           s[i + k] == s[j + k]; k++);
  }
};
```

8.5 Suffix Automaton [016373]

```
struct exSAM {
  const int CNUM = 26;
  // len: maxlength, link: fail link
  // LenSorted: topo order, cnt: occur
  vector<int> len, link, lenSorted, cnt;
  vector<vector<int>> next;
  int total = 0;
  int newnode() {
   return total++;
  void init(int n) { // total number of characters
    len.assign(2 * n, 0); link.assign(2 * n, 0);
lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
    next.assign(2 * n, vector<int>(CNUM));
    newnode(), link[0] = -1;
  int insertSAM(int last, int c) {
    // not exSAM: cur = newnode(), p = Last
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
      next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
```

}

auto add_rep = [&](bool left, int c, int l, int k1,

int k2) {

```
if (len[p] + 1 == len[q]) return link[cur] = q, cur
                                                                 const int L = max(1, 1 - k2), R = min(1 - left, k1)
                                                                if (L > R) return;
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
                                                                if (left) rep.emplace_back(Rep({sft + c - R, sft +
      next[clone][i] = len[next[q][i]] ? next[q][i] :
                                                                     c - L, 1}));
                                                                 else rep.emplace_back(Rep({sft + c - R - l + 1, sft
                                                                      + c - L - l + 1, 1));
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
                                                              for (int cntr = 0; cntr < n; cntr++) {</pre>
    link[link[cur] = clone] = link[q];
                                                                int 1, k1, k2;
    link[q] = clone;
                                                                if (cntr < nu) {</pre>
                                                                  1 = nu - cntr;
    return cur;
  }
                                                                  k1 = get_z(z1, nu - cntr);
  void insert(const string &s) {
                                                                  k2 = get_z(z2, nv + 1 + cntr);
    int cur = 0;
                                                                } else {
                                                                  l = cntr - nu + 1;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
                                                                  k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      if (!nxt) nxt = newnode();
                                                                  k2 = get_z(z4, (cntr - nu) + 1);
      cnt[cur = nxt] += 1;
   }
                                                                 if (k1 + k2 >= 1)
                                                                  add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
  void build() {
                                                              }
    queue<int> q;
    q.push(0);
                                                            8.8 AC Automaton [f529e6]
    while (!q.empty()) {
      int cur = q.front();
                                                            const int SIGMA = 26;
      q.pop();
                                                            struct AC_Automaton {
      for (int i = 0; i < CNUM; ++i)</pre>
                                                              // child: trie, next: automaton
        if (next[cur][i])
                                                              vector<vector<int>> child, next;
          q.push(insertSAM(cur, i));
                                                              vector<int> fail, cnt, ord;
                                                              int total = 0;
    vector<int> lc(total);
                                                              int newnode() {
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
                                                                return total++:
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
                                                              void init(int len) { // len >= 1 + total len
                                                                child.assign(len, vector<int>(26, -1));
        i]]] = i;
                                                                next.assign(len, vector<int>(26, -1));
  void solve() {
                                                                fail.assign(len, -1); cnt.assign(len, 0);
    for (int i = total - 2; i >= 0; --i)
                                                                ord.clear();
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
                                                                newnode();
 }
};
                                                              int input(string &s) {
                                                                int cur = 0;
     Z-value Algorithm [488d87]
8.6
                                                                 for (char c : s) {
                                                                  if (child[cur][c - 'A'] == -1)
  child[cur][c - 'A'] = newnode();
// z[i] = max k s.t. s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
                                                                  cur = child[cur][c - 'A'];
// z[0] = 0
vector<int> z_function(const string &s){
                                                                return cur; // return the end node of string
  int n = s.size();
  vector<int> z(n);
                                                              void make_fl() {
  for(int i = 1, l = 0, r = 0; i < n; i++){</pre>
                                                                queue<int> q;
    if(i \le r) z[i] = min(r - i + 1, z[i - 1]);
                                                                q.push(0), fail[0] = -1;
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
                                                                while(!q.empty()) {
      z[i]++;
                                                                  int R = q.front();
    if(i + z[i] - 1 > r)
                                                                  q.pop(); ord.pb(R);
      1 = i, r = i + z[i] - 1;
                                                                  for (int i = 0; i < SIGMA; i++)</pre>
  }
                                                                    if (child[R][i] != -1) {
  return z;
                                                                       int X = next[R][i] = child[R][i], Z = fail[R
                                                                       while (Z != -1 && child[Z][i] == -1)
8.7
      Main Lorentz [fcfb8f]
                                                                         Z = fail[Z];
struct Rep{ int minl, maxl, len; };
                                                                       fail[X] = Z != -1 ? child[Z][i] : 0;
                                                                       q.push(X);
vector<Rep> rep; // 0-base
// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)
void main_lorentz(const string &s, int sft = 0) {
                                                                     else next[R][i] = R ? next[fail[R]][i] : 0;
  const int n = s.size();
                                                                }
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
                                                              void solve() {
  const string u = s.substr(0, nu), v = s.substr(nu),
                                                                for (int i : ord | views::reverse)
                                                                  if (i) cnt[fail[i]] += cnt[i];
        ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
            ());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
                                                            };
  const auto z1 = z_function(ru), z2 = z_function(v + '
                                                            8.9
                                                                  Palindrome Automaton [8a071b]
                                                            struct PalindromicTree {
             z3 = z_{function}(ru + '#' + rv), z4 =
                 z_function(v);
                                                              struct node {
                                                                int nxt[26], fail, len; // num = depth of fail link
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
                                                                int cnt, num; // cnt = occur, num = #pal_suffix of
                                                                     this node
```

 $node(int 1 = 0) : nxt{}, fail(0), len(1), cnt(0), num$

(0) {}

```
vector<node> st; vector<int> s; int last, n;
  void init() {
    st.clear(); s.clear(); last = 1; n = 0;
    st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
  int getFail(int x) {
    while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x;
 void add(int c) {
   s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
      int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
      st[cur].nxt[c] = now;
      st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
  void dpcnt() {
    for(int i = SZ(st) - 1; i >= 0; i--){
      auto nd = st[i];
      st[nd.fail].cnt += nd.cnt;
  int size() { return (int)st.size() - 2; }
};
```

9 Misc

9.1 Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) ~ f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
   if (n == 1) return 0;
   int l = 0, r = n; bool rv = pred(1, 0);
   while (r - 1 > 1) {
      int m = (1 + r) / 2;
      if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
      else l = m;
   }
   return pred(l, r % n) ? l : r % n;
}
```

9.2 Matroid

 $M=(E,\mathcal{I})$, where $\mathcal{I}\subseteq 2^E$ is nonempty, is a matroid if:

- If $S \in \mathcal{I}$ and $S' \subsetneq S$, then $S' \in \mathcal{I}$.
- For $S_1,S_2\in\mathcal{I}$ s.t. $|S_1|<|S_2|$, there exists $e\in S_2\setminus S_1$ s.t. $S_1\cup\{e\}\in\mathcal{I}$. Matroid intersection:

Start from $S = \emptyset$. In each iteration, let

- $Y_1 = \{x \notin S \mid S \cup \{x\} \in \mathcal{I}_1\}$ • $Y_2 = \{x \notin S \mid S \cup \{x\} \in \mathcal{I}_2\}$
- If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \notin S$,

create edges • $x \to y$ if $S - \{x\} \cup \{y\} \in \mathcal{I}_1$.

• $x \to y \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_1.$ • $y \to x \text{ if } S - \{x\} \cup \{y\} \in \mathcal{I}_2.$

Find a shortest path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x) if $x \notin S$. Find the path with the minimum number of edges among all minimum length paths and alternate it.

9.3 Simulate Annealing [ff826c]

```
ld anneal() {
  mt19937 rnd_engine(seed);
  uniform_real_distribution<ld> rnd(0, 1);
  const ld dT = 0.001;
  // Argument p
  ld S_cur = calc(p), S_best = S_cur;
  for (ld T = 2000; T > eps; T -= dT) {
      // Modify p to p_prime
      const ld S_prime = calc(p_prime);
      const ld delta_c = S_prime - S_cur;
    ld prob = min((ld)1, exp(-delta_c / T));
    if (rnd(rnd_engine) <= prob)</pre>
```

```
S_cur = S_prime, p = p_prime;
if (S_prime < S_best) // find min
    S_best = S_prime, p_best = p_prime;
}
return S_best;
}
9.4 Binary Search On Fraction [f6b9ec]</pre>
```

```
struct Q {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p * d, q + b.q * d};
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N, auto &&pred) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step /= 2 : step *=</pre>
         2);)
      if (Q mid = hi.go(lo, len + step);
           mid.p > N || mid.q > N || dir ^ pred(mid))
        t++;
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

9.5 Min Plus Convolution [09b5c3]

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector
    <int> &b) {
  int n = SZ(a), m = SZ(b);
  vector<int> c(n + m - 1, INF);
  auto dc = [&](auto Y, int 1, int r, int jl, int jr) {
    if (1 > r) return;
    int mid = (1 + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)</pre>
      if (int i = mid - j; i >= 0 && i < n)</pre>
        if (best > a[i] + b[j])
          best = a[i] + b[j], from = j;
    Y(Y, 1, mid - 1, jl, from), Y(Y, mid + 1, r, from,
        ir);
  };
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
```

9.6 **SMAWK** [a2a4ce]

```
// For all 2x2 submatrix:
// If M[1][0] < M[1][1], M[0][0] < M[0][1]
// If M[1][0] == M[1][1], M[0][0] <= M[0][1]
// M[i][ans_i] is the best value in the i-th row
// select(int r, int u, int v) return true if f(r, v)
    is better than f(r, u)
vector<int> smawk(int N, int M, auto &&select) {
  auto dc = [&](auto self, const vector<int> &r, const
      vector<int> &c) {
    if (r.empty()) return vector<int>{};
    const int n = SZ(r); vector<int> ans(n), nr, nc;
    for (int i : c) {
      while (!nc.empty() &&
          select(r[nc.size() - 1], nc.back(), i))
        nc.pop_back();
     if (int(nc.size()) < n) nc.push_back(i);</pre>
    for (int i = 1; i < n; i += 2) nr.push_back(r[i]);</pre>
    const auto na = self(self, nr, nc);
    for (int i = 1; i < n; i += 2) ans[i] = na[i >> 1];
    for (int i = 0, j = 0; i < n; i += 2) {
      ans[i] = nc[j];
      const int end = i + 1 == n ? nc.back() : ans[i +
          1];
      while (nc[j] != end)
        if (select(r[i], ans[i], nc[++j])) ans[i] = nc[
            j];
```

```
return ans;
vector<int> R(N), C(M); iota(iter(R), 0), iota(iter(C)
return dc(dc, R, C);
```

10 Notes

Geometry

Rotation Matrix

$$\begin{pmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{pmatrix}$$

- rotate 90° : $(x,y) \rightarrow (-y,x)$
- rotate -90° : $(x,y) \rightarrow (y,-x)$

Triangles

Side lengths: a,b,c

Semiperimeter:
$$p = \frac{a+b+c}{2}$$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$ Circumradius: $R = \frac{abc}{4A}$

Inradius:
$$r = \frac{A}{n}$$

Length of median (divides triangle into two equal-area triangles): $m_a =$ $\frac{1}{2}\sqrt{2b^2+2c^2-a^2}$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}$$

 $\begin{array}{l} \text{Law of sines: } \frac{\sin\alpha}{a} = \frac{\sin\beta}{b} = \frac{\sin\gamma}{c} = \frac{1}{2R} \\ \text{Law of cosines: } a^2 = b^2 + c^2 - 2bc\cos\alpha \\ \text{Law of tangents: } \frac{a+b}{a-b} = \frac{\tan\frac{\alpha+\beta}{2}}{\tan\frac{\alpha-\beta}{2}} \end{array}$

Ouadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef=ac+bd , and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$. Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \mathrm{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \mathrm{atan2}(y,x) \end{array}$$

Green's Theorem

$$\iint_{D} \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy = \oint_{L^{+}} (Pdx + Qdy)$$
 Area $= \frac{1}{2} \oint_{L} x \ dy - y \ dx$

Circular sector:

$$\begin{split} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r (r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{split}$$

Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1,p_2,p_3 are collinear $\iff p_1^*,p_2^*,p_3^*$ intersect at a point p lies above $l \iff l^*$ lies above p^*
- lower convex hull \leftrightarrow upper envelope

Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x})$$

$$\cosh x = \frac{1}{2}(e^x + e^{-x})$$

$$\sin n\pi = 0$$

$$\cos n\pi = (-1)^n$$

$$\begin{split} \sin(\alpha+\beta) &= \sin\alpha\cos\beta + \cos\alpha\sin\beta \\ \cos(\alpha+\beta) &= \cos\alpha\cos\beta - \sin\alpha\sin\beta \\ \sin(2\alpha) &= 2\cos\alpha\sin\alpha \\ \cos(2\alpha) &= 2\cos^2\alpha - \sin^2\alpha = 2\cos^2\alpha - 1 = 1 - 2\sin^2\alpha \\ \tan(\alpha+\beta) &= \frac{\tan\alpha + \tan\beta}{1 - \tan\alpha\tan\beta} \\ \sin\alpha + \sin\beta &= 2\sin\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \cos\alpha + \cos\beta &= 2\cos\frac{\alpha+\beta}{2}\cos\frac{\alpha-\beta}{2} \\ \sin\alpha\sin\beta &= \frac{1}{2}(\cos(\alpha-\beta) - \cos(\alpha+\beta)) \\ \sin\alpha\cos\beta &= \frac{1}{2}(\sin(\alpha+\beta) + \sin(\alpha-\beta)) \\ \cos\alpha\sin\beta &= \frac{1}{2}(\cos(\alpha-\beta) - \cos(\alpha+\beta)) \\ \cos\alpha\cos\beta &= \frac{1}{2}(\sin(\alpha+\beta) + \sin(\alpha-\beta)) \\ \cos\alpha\cos\beta &= \frac{1}{2}(\cos(\alpha-\beta) + \cos(\alpha+\beta)) \\ (V+W)\tan(\alpha-\beta)/2 &= (V-W)\tan(\alpha+\beta)/2 \\ \text{where } V, W \text{ are lengths of sides opposite angles } \alpha, \beta. \\ &= a\cos x + b\sin x = r\cos(x-\phi) \\ &= a\sin x + b\cos x = r\sin(x+\phi) \end{split}$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

10.3 Calculus

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1-x^{2}}} \qquad \frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^{2}}}$$

$$\frac{d}{dx} \tan x = 1 + \tan^{2}x \qquad \frac{d}{dx} \arctan x = \frac{1}{1+x^{2}}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^{2}}$$

$$\int e^{-x^{2}} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int xe^{ax} = \frac{e^{ax}}{a^{2}} (ax - 1)$$

$$\int \sin^{2}(x) = \frac{x}{2} - \frac{1}{4} \sin 2x \qquad \int \sin^{3}x = \frac{1}{12} \cos 3x - \frac{3}{4} \cos x$$

$$\int \cos^{2}(x) = \frac{x}{2} + \frac{1}{4} \sin 2x \qquad \int \cos^{3}x = \frac{1}{12} \sin 3x + \frac{3}{4} \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int xe^{x} = e^{x}(x - 1) \qquad \int x^{2}e^{x} = e^{x}(x^{2} - 2x + 2)$$

$$\int x^{2} \sin x = 2x \sin x - (x^{2} - 2) \cos x$$

$$\int x^{2} \cos x = 2x \cos x + (x^{2} - 2) \sin x$$

$$\int e^{x} \sin x = \frac{1}{2}e^{x}(\sin x - \cos x)$$

$$\int e^{x} \cos x = \frac{1}{2}e^{x}(\sin x - \cos x)$$

$$\int xe^{x} \sin x = \frac{1}{2}e^{x}(x \sin x - x \cos x + \cos x)$$

$$\int xe^{x} \cos x = \frac{1}{2}e^{x}(x \sin x - x \cos x - \sin x)$$

10.4 Sum & Series

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

10.5 Misc

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- BEST theorem: the number of eulerian circuits in a directed graph is $|\det(L_{ww})| \cdot \prod_{v \in V} (\deg(v) - 1)!$.
- Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

- · Cayley's Formula
 - Given a degree sequence d_1, d_2, \ldots, d_n for each $\emph{labeled}$ vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.
- Erdős-Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only

if
$$d_1+\cdots+d_n$$
 is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds

for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n

is bigraphic if and only if
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

holds for every $1 \le k \le n$.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1,b_1),\ldots,(a_n,b_n)$ of nonnegative integer pairs with $a_1\geq$

$$\cdots \geq a_n$$
 is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k-1)$

$$1) + \sum_{i=k+1}^n \min(b_i, k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have $A=\#\{\mbox{lattice points in the interior}\}+\frac{\#\{\mbox{lattice points on the boundary}\}}{2}-1.$

- Möbius inversion formula
 - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
 - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- - A portion of a sphere cut off by a plane.
 - r: sphere radius, a: radius of the base of the cap, h: height of the cap,
 - Volume = $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$ $\cos \theta)^2/3$.
 - Area = $2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$.
- · Lagrange multiplier
 - Optimize $f(x_1,\ldots,x_n)$ when k constraints $g_i(x_1,\ldots,x_n)=0$.
 - Lagrangian function
 - Lagrangian random $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k)=f(x_1,\ldots,x_n)-\sum_{i=1}^k\lambda_ig_i(x_1,\ldots,x_n).$ The solution corresponding to the original constrained optimization
 - is always a saddle point of the Lagrangian function.
- Nearest points of two skew lines

– Line 1 :
$$oldsymbol{v}_1 = oldsymbol{p}_1 + t_1 oldsymbol{d}_1$$

- Line 2 :
$${m v}_2 = {m p}_2 + t_2 {m d}_2$$

- $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$

-
$$c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1$$

$$egin{align*} & \mathbf{a} = \mathbf{d}_1 imes \mathbf{d}_2 \\ & - \mathbf{n}_1 = \mathbf{d}_1 imes \mathbf{n} \\ & - \mathbf{n}_2 = \mathbf{d}_2 imes \mathbf{n} \\ & - \mathbf{c}_1 = \mathbf{p}_1 + rac{(\mathbf{p}_2 - \mathbf{p}_1) \cdot \mathbf{n}_2}{\mathbf{d}_1 \cdot \mathbf{n}_2} \mathbf{d}_1 \\ & - \mathbf{c}_2 = \mathbf{p}_2 + rac{(\mathbf{p}_1 - \mathbf{p}_2) \cdot \mathbf{n}_1}{\mathbf{d}_2 \cdot \mathbf{n}_1} \mathbf{d}_2 \\ & - \mathbf{n}_1 = \mathbf{n}_1 \mathbf{n}_2 \mathbf{n}_2 \mathbf{n}_1 \mathbf{n}_2 \mathbf$$

• Bernoulli numbers
$$B_0-1, B_1^{\pm}=\pm\frac{1}{2}, B_2=\frac{1}{6}, B_3=0$$

$$\sum_{j=0}^m {m+1 \choose j} B_j = 0 \text{, EGF is } B(x) = \frac{x}{e^x-1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

- Stirling numbers of the second kind Partitions of \boldsymbol{n} distinct elements into exactly k groups.

$$\begin{array}{l} S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) = \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n = \sum_{i=0}^n S(n,i)(x)_i \\ \bullet \text{ Pentagonal number theorem} \end{array}$$

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} {n+1 \choose j} (k+1-j)^{n}$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{k=0}^{k} (-1)^{j} {n+1 \choose k} (k+1-j)^{n}$$

10.6 Number

· Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039, 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

• Number of paritions of n:

Maximum number of divisors:

d(i) 12 32 240 1344 6720 26880 103680

n | 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 (2n) 2 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

• Number of ways to partition a set of n labeled elements: $n \mid 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \ 13$

$$\frac{n}{B_n}$$
 2 5 15 52 203 877 4140 21147 115975 7e5 4e6 3e7

• Fibonacci numbers: $\frac{n}{F_n}$ 1 1 2 3 4 5 31 45 88 1346269 1e9 1e18

