Contents 5.11 Minimum Steiner Tree . . . 14 const int N = 1 << 20; 5.12 Count Cycles 15 static char buf[N]; static char *p = buf , *end = buf; Basic 1.1 .vimrc if(p == end) { 6.1 Extended Euclidean Algo-1.2 Fast IO rithm 15 Random 6.2 Floor & Ceil 15 1.4 PBDS Tree p = buf;1.5 Pragma 6.3 Legendre 15 } 1.6 SVG Writer 6.4 Simplex 15 return *p++; 6.5 Simplex Construction . . . 16 2 Data Structure 6.6 DiscreteLog 16 2.1 Heavy-Light Decomposition 6.7 Miller Rabin & Pollard Rho . 16 2.2 Link Cut Tree 6.8 XOR Basis 16 2.3 Treap struct Writer { 2.4 KD Tree 6.9 Linear Equation 16 2.5 Leftist Tree 6.10 Chinese Remainder Theorem 17 Convex 1D/1D . . 6.11 Sqrt Decomposition 17 2.7 Dynamic Convex Hull 6.12 Floor Sum 17 (); } Flow & Matching 7 Polynomial 3.1 Dinic 3.2 Bounded Flow 7.1 FWHT 17 3.3 MCMF 7.2 FFT 17 3.4 Min Cost Circulation 7.3 NTT 18 3.5 Gomory Hu 7.4 Polynomial Operation . . . 18 3.6 Stoer Wagner Algorithm . . 7.5 Generating Function 20 3.7 Bipartite Matching Ordinary Generating Func-3.8 Kuhn Munkres Algorithm tion 20 Exponential Generating 3.9 Max Simple Graph Matching 6 3.10 Flow Model x % 10; Function 20 7.6 Bostan Mori 20 Geometry size += ptr; 4.1 Geometry Template 4.2 Polar Angle Comparator . . 8 String Minkowski Sum 8.1 KMP Algorithm 21 Intersection of Circle and 8.2 Manacher Algorithm 21 Convex Polygon 8 8.3 Lyndon Factorization 21 Intersection of Circles . . . 8.4 Suffix Array 21 Tangent Line of Circles . . . 8.5 Suffix Automaton 21 Intersection of Line and 4.7 8.6 Z-value Algorithm 22 Convex Polygon 8 8.7 Main Lorentz 22 Intersection of Line and Circle 8.8 AC Automaton 22 4.9 Point in Circle 8 8.9 Palindrome Automaton . . 22 4.10 Point in Convex 8.10 Palindrome Partition 23 4.11 Half Plane Intersection . . 4.12 HPI General Line 4.13 Minimum Enclosing Circle . 9.1 Cyclic Ternary Search . . . 23 9.2 Matroid 23 9.3 Simulate Annealing 23 // .order_of_key(x) 4.16 Delaunay Triangulation . . 10 4.17 Voronoi Diagram 10 9.4 Binary Search On Fraction . 23 4.18 Polygon Union 9.5 Min Plus Convolution . . . 23 4.19 Tangent Point to Convex Hull 11 9.6 SMAWK 23 4.20 Heart 11 9.7 Golden Ratio Search 24 4.21 Rotating Sweep Line 11 10 Notes 10.1 Geometry 4.24 Calculate Points in Triangle 11 Rotation Matrix 24 Triangles 24 Graph Quadrilaterals 24 Spherical coordinates . . . 24 Green's Theorem 24 5.4 Dominator Tree 12 Point-Line Duality 24 10.2 Trigonometry 24 5.6 Fast DMST 13 10.3 Calculus 24 10.4 Sum & Series 25 5.9 Number of Maximal Clique 14 5.10 Minimum Mean Cycle . . . 14 10.5 Misc 25 10.6 Number 25 Basic Default code: Basic 9c8f02, Debug 28c438) { 1.1 .vimrc [c107f4] se nu rnu bs=2 sw=4 ts=4 hls ls=2 si acd bo=all mouse=a map <F9> :w<bar>!g++ "%" -o %:r -std=c++20 -Wall -Wextra -Wshadow -O2 -Dzisk -g -fsanitize=address, undefined<CR> map <F8> :!./%:r<CR> inoremap {<CR> {<CR>}<ESC>ko ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space :]' \| md5sum \| cut -c-6 inoremap fj <ESC> vnoremap fj <ESC> -D_GLIBCXX_ASSERTIONS, -D_GLIBCXX_DEBUG

1.2 Fast IO [4f6f0e]

char readchar() {

```
if((end = buf + fread(buf , 1 , N , stdin)) == buf)
           return EOF;
const int buf_size = 524288;
  char buf[buf_size]; int size = 0, ret;
  void flush() { ret = write(1, buf, size); size = 0; }
         _flush(<mark>int</mark> sz) {        <mark>if</mark> (sz + size > buf_size) flush
  void write_char(char c) { _flush(1); buf[size++] = c;
  void write_int(int x) {
    const int len = 20;
     if (x == 0) buf[size + (ptr++)] = '0';
     else for (; x; x /= 10) buf[size + (ptr++)] = '0' +
     reverse(buf + size, buf + size + ptr);
}; // remember to call flush
1.3 Random [4cf9ed]
mt19937 rng(chrono::system_clock::now().
     time_since_epoch().count());
1.4 PBDS Tree [9e57e3]
#include <bits/extc++.h>
using namespace __gnu_pbds;
using Tree = tree<int, null_type, less<>, rb_tree_tag,
     tree_order_statistics_node_update>;
// .find_by_order(x)
1.5 Pragma [6006f6]
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
 _builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
1.6 SVG Writer [7adcc8]
  void p(string_view s) { o << s; }</pre>
  void p(string_view s, auto v, auto... vs) {
  auto i = s.find('$');
     o << s.substr(0, i) << v, p(s.substr(i + 1), vs...)
  ofstream o; string c = "red";
public: // SVG svg("test.svg", 0, 0, 100, 100)
    SVG(auto f, auto x1, auto y1, auto x2, auto y2) : o(f
     p("<svg xmlns='http://www.w3.org/2000/svg' "
        'viewBox='$ $ $ $'>\n"
       "<style>*{stroke-width:0.5%;}</style>\n",
  x1, -y2, x2 - x1, y2 - y1); }
~SVG() { p("</svg>\n"); }
  void color(string nc) { c = nc; }
  void line(auto x1, auto y1, auto x2, auto y2) {
  p("<line x1='$' y1='$' x2='$' y2='$' stroke='$'/>\n
       x1, -y1, x2, -y2, c); }
  void circle(auto x, auto y, auto r) {
  p("<circle cx='$' cy='$' r='$' stroke='$' "
    "fill='none'/>\n", x, -y, r, c); }
  void text(auto x, auto y, string s, int w = 12) {
  p("<text x='$' y='$' font-size='$px'>$</text>\n",
       x, -y, w, s); }
};
```

Data Structure

2.1 Heavy-Light Decomposition [f2dbca]

```
struct HLD{ // 1-based
 int n, ts = 0; // ord is 1-based
  vector<vector<int>> g;
  vector<int> par, top, down, ord, dpt, sub;
explicit HLD(int _n): n(_n), g(n + 1),
  par(n + 1), top(n + 1), down(n + 1), ord(n + 1), dpt(n + 1), sub(n + 1) {}
  void add_edge(int u, int v){ g[u].pb(v); g[v].pb(u);
  void dfs(int now, int p){
    par[now] = p; sub[now] = 1;
    for(int i : g[now]){
      if(i == p) continue;
      dpt[i] = dpt[now] + 1;
      dfs(i, now);
      sub[now] += sub[i];
      if(sub[i] > sub[down[now]]) down[now] = i;
    }
  void cut(int now, int t){
    top[now] = t; ord[now] = ++ts;
    if(!down[now]) return;
    cut(down[now], t);
    for(int i : g[now]){
      if(i != par[now] && i != down[now])
        cut(i, i);
    }
  void build(){ dfs(1, 1), cut(1, 1); }
  int query(int a, int b){
    int ta = top[a], tb = top[b];
    while(ta != tb){
      if(dpt[ta] > dpt[tb]) swap(ta, tb), swap(a, b);
      // ord[tb], ord[b]
      tb = top[b = par[tb]];
    if(ord[a] > ord[b]) swap(a, b);
    // ord[a], ord[b]
    return a; // Lca
};
```

2.2 Link Cut Tree [502ab1]

```
// 1-based
// == 43515a ==
template <typename Val, typename SVal> struct LCT {
  struct node {
    int pa, ch[2]; bool rev; int size;
    Val v, sum, rsum; SVal sv, sub, vir;
    node(): pa{0}, ch{0, 0}, rev{false}, size{1}, v{},
      sum{}, rsum{}, sv{}, sub{}, vir{} {}
#define cur o[u]
#define lc cur.ch[0]
#define rc cur.ch[1]
  vector<node> o;
  bool is_root(int u) const {
    return o[cur.pa].ch[0]!=u && o[cur.pa].ch[1]!=u; }
  bool is rch(int u) const {
    return o[cur.pa].ch[1] == u && !is_root(u); }
  void down(int u) {
    for (int c : {lc, rc}) if (c) {
      if (cur.rev) set_rev(c);
    cur.rev = false;
  void up(int u) {
    cur.sum = o[lc].sum + cur.v + o[rc].sum;
    cur.rsum = o[rc].rsum + cur.v + o[lc].rsum;
    cur.sub = cur.vir + o[lc].sub + o[rc].sub + cur.sv;
    cur.size = o[lc].size + o[rc].size + 1;
  void set_rev(int u) {
    swap(lc, rc), swap(cur.sum, cur.rsum);
    cur.rev ^= 1;
// == 3a186b ==
  void rotate(int u) {
```

```
int f = cur.pa, g = o[f].pa, l = is_rch(u);
if (cur.ch[1 ^ 1]) o[cur.ch[1 ^ 1]].pa = f;
    if (not is_root(f)) o[g].ch[is_rch(f)] = u;
    o[f].ch[l] = cur.ch[l ^ 1], cur.ch[l ^ 1] = f;
    cur.pa = g, o[f].pa = u; up(f);
  vector<int> stk;
  void splay(int u) {
    stk.clear(); stk.pb(u);
    while (not is_root(stk.back()))
      stk.push_back(o[stk.back()].pa);
    while (not stk.empty())
      down(stk.back()), stk.pop_back();
    for (int f = cur.pa; not is_root(u); f = cur.pa) {
      if (!is_root(f))
        rotate(is_rch(u) == is_rch(f) ? f : u);
      rotate(u);
    }
    up(u);
  }
  void access(int x) {
    for (int u = x, last = 0; u; u = cur.pa) {
      splay(u);
      cur.vir = cur.vir + o[rc].sub - o[last].sub;
      rc = last; up(last = u);
    splay(x);
  int find_root(int u) {
    int la = 0;
    for (access(u); u; u = lc) down(la = u);
    return la:
  void split(int x, int y) { chroot(x); access(y); }
  void chroot(int u) { access(u); set_rev(u); }
// == a238c2 ==
  LCT(int n = 0) : o(n + 1) { o[0].size = 0; }
  void set_val(int u, const Val &v) {
    splay(u); cur.v = v; up(u); }
  void set_sval(int u, const SVal &v) {
    access(u); cur.sv = v; up(u); }
  Val query(int x, int y) {
    split(x, y); return o[y].sum; }
  SVal subtree(int p, int u) {
    chroot(p); access(u); return cur.vir + cur.sv; }
  bool connected(int u, int v) {
    return find_root(u) == find_root(v); }
  void link(int x, int y) {
    chroot(x); access(y);
    o[y].vir = o[y].vir + o[x].sub;
    up(o[x].pa = y);
  void cut(int x, int y) {
    split(x, y); o[y].ch[0] = o[x].pa = 0; up(y); }
#undef cur
#undef lc
#undef ro
};
2.3 Treap [2ac37e]
mt19937 rng(880301);
// == fb4359 ==
struct node {
  11 data; int sz;
  node *1, *r;
  node(ll k = 0) : data(k), sz(1), l(0), r(0) {}
  void up() {
    if (1) sz += 1->sz;
    if (r) sz += r->sz;
  void down() {}
};
node pool[1000010]; int pool_cnt = 0;
node *newnode(l1 k){ return &(pool[pool_cnt++] = node(k
    )); }
int sz(node *a) { return a ? a->sz : 0; }
node *merge(node *a, node *b) {
  if (!a | | !b) return a ? a : b;
  if (int(rng() % (sz(a) + sz(b))) < sz(a))</pre>
    return a->down(), a->r = merge(a->r, b), a->up(),
```

a;

```
return b->down(), b->l = merge(a, b->l), b->up(), b;
                                                                   return m;
// a: key <= k, b: key > k
void split(node *o, node *&a, node *&b, ll k) {
                                                                 bool bound(const point &q, int o, long long d) {
  if (!o) return a = b = 0, void();
                                                                   double ds = sqrt(d + 1.0);
  o->down();
                                                                    if (q.x < x1[o] - ds || q.x > xr[o] + ds ||
                                                                        q.y < y1[o] - ds || q.y > yr[o] + ds)
  if (o->data <= k)
    a = o, split(o->r, a->r, b, k), <math>a->up();
                                                                      return false;
  else b = o, split(o->1, a, b->1, k), b->up();
                                                                   return true;
// a: size k, b: size n - k
                                                                 long long dist(const point &a, const point &b) {
void split2(node *o, node *&a, node *&b, int k) {
                                                                   return (a.x - b.x) * 111 * (a.x - b.x) +
                                                                     (a.y - b.y) * 111 * (a.y - b.y);
  if (sz(o) <= k) return a = o, b = 0, void();</pre>
  o->down();
                                                                 void dfs(
  if (sz(o->1) + 1 <= k)
                                                                     const point &q, long long &d, int o, int dep = 0)
    a = o, split2(o->r, a->r, b, k - <math>sz(o->l) - 1);
  else b = o, split2(o->1, a, b->1, k);
  o->up();
                                                                   if (!bound(q, o, d)) return;
}
                                                                   long long cd = dist(p[o], q);
// == e9f4d8 ==
                                                                    if (cd != 0) d = min(d, cd);
node *kth(node *o, ll k) { // 1-based
                                                                   if ((dep & 1) && q.x < p[o].x ||</pre>
  if (k <= sz(o->1)) return kth(o->1, k);
                                                                        !(dep & 1) && q.y < p[o].y) {
                                                                      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  if (k == sz(o->1) + 1) return o;
  return kth(o\rightarrow r, k - sz(o\rightarrow 1) - 1);
                                                                     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                   } else {
int Rank(node *o, ll key) { // num of key < key</pre>
                                                                     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  if (!o) return 0;
                                                                     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  if (o->data < key)</pre>
                                                                   }
    return sz(o->1) + 1 + Rank(o->r, key);
                                                                 }
  else return Rank(o->1, key);
                                                                 void init(const vector<point> &v) {
                                                                   for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
bool erase(node *&o, ll k) {
                                                                   root = build(0, v.size());
  if (!o) return 0;
  if (o->data == k) {
                                                                 long long nearest(const point &q) {
    node *t = o;
                                                                   long long res = 1e18;
    o->down(), o = merge(o->1, o->r);
                                                                   dfs(q, res, root);
    return 1;
                                                                   return res:
  node *&t = k < o->data ? o->l : o->r;
                                                               } // namespace kdt
  return erase(t, k) ? o->up(), 1 : 0;
                                                               2.5 Leftist Tree [e91538]
void insert(node *&o, ll k) {
                                                               struct node {
  node *a, *b;
                                                                 11 v, data, sz, sum;
node *1, *r;
  split(o, a, b, k),
    o = merge(a, merge(new node(k), b));
                                                                 node(ll k)
                                                                    : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
tuple<node*, node*, node*> interval(node *&o, int 1,
  int r) { // 1-based
node *a, *b, *c; // b: [l, r]
                                                               11 sz(node *p) { return p ? p->sz : 0; }
                                                               11 V(node *p) { return p ? p->v : -1; }
  split2(o, a, b, l - 1), split2(b, b, c, r - l + 1);
                                                               11 sum(node *p) { return p ? p->sum : 0; }
  return make_tuple(a, b, c);
                                                               node *merge(node *a, node *b) {
}
                                                                 if (!a || !b) return a ? a : b;
2.4 KD Tree [375ca2]
                                                                 if (a->data < b->data) swap(a, b);
                                                                 a->r = merge(a->r, b);
namespace kdt {
                                                                 if (V(a\rightarrow r) \rightarrow V(a\rightarrow l)) swap(a\rightarrow r, a\rightarrow l);
  int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
                                                                 a -> v = V(a -> r) + 1, a -> sz = sz(a -> 1) + sz(a -> r) + 1;
  yl[maxn], yr[maxn];
                                                                 a\rightarrow sum = sum(a\rightarrow 1) + sum(a\rightarrow r) + a\rightarrow data;
  point p[maxn];
                                                                 return a;
  int build(int 1, int r, int dep = 0) {
    if (1 == r) return -1;
                                                               void pop(node *&o) {
    function<bool(const point &, const point &)> f =
                                                                 node *tmp = o;
      [dep](const point &a, const point &b) {
                                                                 o = merge(o->1, o->r);
        if (dep & 1) return a.x < b.x;
                                                                 delete tmp;
        else return a.y < b.y;</pre>
                                                               }
      };
    int m = (1 + r) >> 1;
                                                               2.6 Convex 1D/1D [a449dd]
    nth_element(p + 1, p + m, p + r, f);
    x1[m] = xr[m] = p[m].x;
                                                               template < class T>
    yl[m] = yr[m] = p[m].y;
                                                               struct DynamicHull {
    lc[m] = build(1, m, dep + 1);
                                                                 struct seg { int x, 1, r; };
                                                                 T f; int C; deque<seg> dq; // range: 1~C explicit DynamicHull(T _f, int _C): f(_f), C(_C) {}
    if (~lc[m]) {
      xl[m] = min(xl[m], xl[lc[m]]);
                                                                 // max t s.t. f(x, t) \Rightarrow f(y, t), x < y, maintain max
      xr[m] = max(xr[m], xr[lc[m]]);
      yl[m] = min(yl[m], yl[lc[m]]);
                                                                 int intersect(int x, int y) {
      yr[m] = max(yr[m], yr[lc[m]]);
                                                                   int 1 = 0, r = C + 1;
                                                                   while (1 + 1 < r) {
    rc[m] = build(m + 1, r, dep + 1);
                                                                     int mid = (1 + r) / 2;
                                                                      if (f(x, mid) >= f(y, mid)) 1 = mid;
    if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
                                                                     else r = mid;
      xr[m] = max(xr[m], xr[rc[m]]);
      yl[m] = min(yl[m], yl[rc[m]]);
yr[m] = max(yr[m], yr[rc[m]]);
                                                                   return 1;
```

```
void push_back(int x) {
    for (int i; !dq.empty() &&
        (i = dq.back().1, f(dq.back().x, i) < f(x, i));
      dq.pop_back();
    if (dq.empty()) return dq.pb(seg({x, 1, C})), void
        ();
    dq.back().r = intersect(dq.back().x, x);
    if (dq.back().r + 1 <= C) dq.pb(seg({x, dq.back().r</pre>
         + 1, C}));
  int query(int x) {
    while (dq.front().r < x) dq.pop_front();</pre>
    return dq.front().x;
 }
};
```

2.7 Dynamic Convex Hull [b45ebc]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable 11 a, b, p;
  bool operator<(const Line &rhs) const { return a <</pre>
      rhs.a; }
 bool operator<(11 x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return 0; }
    if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf
    else x -> p = iceil(y -> b - x -> b, x -> a - y -> a);
    return x->p >= y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() \&\& isect(--x, y)) isect(x, y =
    while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  11 query(11 x) {
    auto 1 = *lower_bound(x);
    return 1.a * x + 1.b;
 }
};
```

Flow & Matching 3

3.1 Dinic [801a71]

```
struct Dinic { // 0-based, O(V^2E), unit flow: O(min(V
    ^{2/3}E, E^{3/2}), bipartite matching: O(sqrt(V)E)
  struct edge {
   11 to, cap, flow, rev;
  int n, s, t;
  vector<vector<edge>> g;
 vector<int> dis, ind;
  void init(int _n) {
   n = _n;
   g.assign(n, vector<edge>());
  void reset() {
   for (int i = 0; i < n; ++i)</pre>
     for (auto &j : g[i]) j.flow = 0;
  void add_edge(int u, int v, ll cap) {
   g[u].pb(edge\{v, cap, 0, SZ(g[v])\});
   g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
   //change g[v] to cap for undirected graphs
  bool bfs() {
   dis.assign(n, -1);
   queue<int> q;
   q.push(s), dis[s] = 0;
   while (!q.empty()) {
     int cur = q.front(); q.pop();
```

```
4
         if (dis[e.to] == -1 && e.flow != e.cap) {
           q.push(e.to);
           dis[e.to] = dis[cur] + 1;
      }
    }
    return dis[t] != -1;
  11 dfs(int u, 11 cap) {
    if (u == t || !cap) return cap;
    for (int &i = ind[u]; i < SZ(g[u]); ++i) {</pre>
      edge &e = g[u][i];
      if (dis[e.to] == dis[u] + 1 && e.flow != e.cap) {
        11 df = dfs(e.to, min(e.cap - e.flow, cap));
         if (df) {
           e.flow += df;
           g[e.to][e.rev].flow -= df;
           return df:
        }
      }
    }
    dis[u] = -1;
    return 0;
  il maxflow(int _s, int _t) {
    s = _s; t = _t;
    11 f\overline{1}ow = 0, df;
    while (bfs()) {
      ind.assign(n, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow:
  }
};
      Bounded Flow [758826]
struct BoundedFlow : Dinic {
  vector<ll> tot;
  void init(int _n) {
    Dinic::init(_n + 2);
    tot.assign(n, 0);
  void add_edge(int u, int v, ll lcap, ll rcap) {
    tot[u] -= lcap, tot[v] += lcap;
    g[u].pb(edge{v, rcap, lcap, SZ(g[v])});
g[v].pb(edge{u, 0, 0, SZ(g[u]) - 1});
```

3.2

```
bool feasible() {
    11 \text{ sum = 0;}
    int vs = n - 2, vt = n - 1;
    for(int i = 0; i < n - 2; ++i)</pre>
      if(tot[i] > 0)
         add_edge(vs, i, 0, tot[i]), sum += tot[i];
       else if(tot[i] < 0) add_edge(i, vt, 0, -tot[i]);</pre>
    if(sum != maxflow(vs, vt)) sum = -1;
    for(int i = 0; i < n - 2; i++)</pre>
      if(tot[i] > 0)
         g[vs].pop_back(), g[i].pop_back();
       else if(tot[i] < 0)</pre>
         g[i].pop_back(), g[vt].pop_back();
    return sum != -1;
  11 boundedflow(int _s, int _t) {
    add_edge(_t, _s, 0, INF);
if(!feasible()) return -1;
    11 x = g[_t].back().flow;
    g[_t].pop_back(), g[_s].pop_back();
    return x - maxflow(_t, _s); // min
    //return x + maxflow(_s, _t); // max
  }
};
```

3.3 MCMF [671e14]

```
struct MCMF { // 0-base
 struct Edge {
   11 from, to, cap, flow, cost, rev;
 int n, s, t;
 vector<vector<Edge>> g;
 vector<Edge*> past;
 vector<ll> dis, up, pot;
```

```
explicit MCMF(int
                      _n): n(_n), g(n), past(n), dis(n),
      up(n), pot(n) \frac{1}{\{}
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(Edge{a, b, cap, 0, cost, SZ(g[b])});
g[b].pb(Edge{b, a, 0, 0, -cost, SZ(g[a]) - 1});
  bool BellmanFord() {
    vector<bool> inq(n);
    fill(iter(dis), INF);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u]) {
        11 d2 = dis[u] + e.cost + pot[u] - pot[e.to];
        relax(e.to, d2, min(up[u], e.cap - e.flow), &e)
      }
    }
    return dis[t] != INF;
  pair<ll, ll> solve(int _s, int _t, bool neg = true) {
    s = _s, t = _t; 11 flow = 0, cost = 0;
    if (neg) BellmanFord(), pot = dis;
    for (; BellmanFord(); pot = dis) {
      for (int i = 0; i < n; ++i)</pre>
        if (dis[i] != INF) dis[i] += pot[i] - pot[s];
      flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], g[e.to][e.rev].flow -= up[t];
    }
    return {flow, cost};
  }
};
    Min Cost Circulation [47cf18]
3.4
struct MinCostCirculation { // 0-based, O(VE * ElogC)
  struct edge {
   11 from, to, cap, fcap, flow, cost, rev;
  int n;
  vector<edge*> past;
  vector<vector<edge>> g;
  vector<ll> dis;
  void BellmanFord(int s) {
    vector<int> inq(n);
    dis.assign(n, INF);
    queue<int> q;
    auto relax = [&](int u, ll d, edge *e) {
      if (dis[u] > d) {
        dis[u] = d, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, 0);
    while (!q.empty()) {
      int u = q.front();
      q.pop(), inq[u] = 0;
      for (auto &e : g[u])
        if (e.cap > e.flow)
          relax(e.to, dis[u] + e.cost, &e);
    }
  void try_edge(edge &cur) {
    if (cur.cap > cur.flow) return ++cur.cap, void();
    BellmanFord(cur.to);
    if (dis[cur.from] + cur.cost < 0) {</pre>
      ++cur.flow, --g[cur.to][cur.rev].flow;
for (int i = cur.from; past[i]; i = past[i]->from
          ) {
        auto &e = *past[i];
        ++e.flow, --g[e.to][e.rev].flow;
```

```
++cur.cap:
  }
  void solve(int mxlg) { // mxlg >= log(max cap)
    for (int b = mxlg; b >= 0; --b) {
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : g[i])
           e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : g[i])
           if (e.fcap >> b & 1)
             try_edge(e);
    }
  void init(int _n) {
    n = _n;
    past.assign(n, nullptr);
    g.assign(n, vector<edge>());
  void add_edge(ll a, ll b, ll cap, ll cost) {
    g[a].pb(edge{a, b, 0, cap, 0, cost, SZ(g[b]) + (a)}
         == b));
    g[b].pb(edge{b, a, 0, 0, 0, -cost, SZ(g[a]) - 1});
  }
};
3.5
      Gomory Hu [82d968]
void GomoryHu(Dinic &flow) { // 0-based
  int n = flow.n:
  vector<int> par(n);
  for (int i = 1; i < n; ++i) {</pre>
    flow.reset();
    add_edge(i, par[i], flow.maxflow(i, par[i]));
for (int j = i + 1; j < n; ++j)
      if (par[j] == par[i] && ~flow.dis[j])
         par[j] = i;
}
      Stoer Wagner Algorithm [a9917b]
3.6
struct StoerWagner { // 0-based, 0(V^3)
  int n;
  vector<int> vis, del;
  vector<ll> wei;
  vector<vector<ll>> edge;
  void init(int _n) {
    n = _n;
    del.assign(n, 0);
    edge.assign(n, vector<ll>(n));
  void add_edge(int u, int v, ll w) {
    edge[u][v] += w, edge[v][u] += w;
  void search(int &s, int &t) {
    vis.assign(n, 0); wei.assign(n, 0);
    s = t = -1:
    while (1) {
      11 mx = -1, cur = 0;
for (int i = 0; i < n; ++i)</pre>
         if (!del[i] && !vis[i] && mx < wei[i])</pre>
           cur = i, mx = wei[i];
      if (mx == -1) break;
      vis[cur] = 1, s = t, t = cur;
      for (int i = 0; i < n; ++i)</pre>
         if (!vis[i] && !del[i]) wei[i] += edge[cur][i];
    }
  11 solve() {
    11 \text{ ret} = INF;
    for (int i = 0, x=0, y=0; i < n-1; ++i) {</pre>
      search(x, y), ret = min(ret, wei[y]), del[y] = 1;
for (int j = 0; j < n; ++j)</pre>
```

3.7 Bipartite Matching [5bb9be]

return ret;

};

// O(E sqrt(V)), O(E log V) for random sparse graphs

edge[x][j] = (edge[j][x] += edge[y][j]);

```
struct BipartiteMatching { // 0-based
 int nl. nr:
  vector<int> mx, my, dis, cur;
  vector<vector<int>> g;
  bool dfs(int u) {
    for (int &i = cur[u]; i < SZ(g[u]); ++i) {</pre>
      int e = g[u][i];
      if (!\sim my[e] \mid | (dis[my[e]] == dis[u] + 1 && dfs(
          my[e])))
        return mx[my[e] = u] = e, 1;
    dis[u] = -1;
    return 0;
  bool bfs() {
   int ret = 0;
    queue<int> q;
    dis.assign(nl, -1);
    for (int i = 0; i < nl; ++i)</pre>
      if (!~mx[i]) q.push(i), dis[i] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (int e : g[u])
        if (!~my[e]) ret = 1;
        else if (!~dis[my[e]]) {
          q.push(my[e]);
          dis[my[e]] = dis[u] + 1;
        }
    return ret;
  int matching() {
   int ret = 0;
    mx.assign(nl, -1); my.assign(nr, -1);
    while (bfs()) {
     cur.assign(nl, 0);
      for (int i = 0; i < nl; ++i)</pre>
        if (!~mx[i] && dfs(i)) ++ret;
    return ret;
  void add_edge(int s, int t) { g[s].pb(t); }
  void init(int _nl, int _nr) {
   nl = _nl, nr = _nr;
    g.assign(nl, vector<int>());
};
     Kuhn Munkres Algorithm [683e0a]
struct KM { // 0-based, maximum matching, O(V^3)
```

```
int n, ql, qr;
vector<vector<ll>> w;
vector<ll> hl, hr, slk;
vector<int> fl, fr, pre, qu, vl, vr;
void init(int _n) {
 // -INF for perfect matching
 w.assign(n, vector<ll>(n, 0));
 pre.assign(n, 0);
 qu.assign(n, 0);
void add_edge(int a, int b, ll wei) {
 w[a][b] = wei;
bool check(int x) {
 if (vl[x] = 1, \sim fl[x])
   return (vr[qu[qr++] = fl[x]] = 1);
 while (\sim x) swap(x, fr[fl[x] = pre[x]]);
 return 0;
}
void bfs(int s) {
  slk.assign(n, INF); vl.assign(n, 0); vr.assign(n,
  ql = qr = 0, qu[qr++] = s, vr[s] = 1;
  for (11 d;;) {
   while (ql < qr)</pre>
      for (int x = 0, y = qu[ql++]; x < n; ++x)
        if (!v1[x] \&\& s1k[x] >= (d = h1[x] + hr[y] -
            w[x][y])) {
          if (pre[x] = y, d) slk[x] = d;
          else if (!check(x)) return;
```

```
d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] && d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {
        if (vl[x]) hl[x] += d;
        else slk[x] -= d;
        if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!v1[x] && !s1k[x] && !check(x)) return;
    }
  11 solve() {
    fl.assign(n, -1); fr.assign(n, -1); hl.assign(n, 0)
    ; hr.assign(n, 0);
for (int i = 0; i < n; ++i)</pre>
      hl[i] = *max_element(iter(w[i]));
    for (int i = 0; i < n; ++i) bfs(i);</pre>
    11 \text{ res} = 0;
    for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
    return res;
};
3.9
     Max Simple Graph Matching [907d7c]
struct Matching { // 0-based, 0(V^3)
  queue<int> q; int n;
  vector<int> fa, s, vis, pre, match;
  vector<vector<int>> g;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
```

```
int LCA(int x, int y) {
  static int tk = 0; tk++; x = Find(x); y = Find(y);
  for (;; swap(x, y)) if (x != n) {
    if (vis[x] == tk) return x;
    vis[x] = tk;
    x = Find(pre[match[x]]);
}
void Blossom(int x, int y, int 1) {
  for (; Find(x) != 1; x = pre[y]) {
    pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
    for (int z: {x, y}) if (fa[z] == z) fa[z] = 1;
bool Bfs(int r) {
  iota(iter(fa), 0); fill(iter(s), -1);
  q = queue<int>(); q.push(r); s[r] = 0;
for (; !q.empty(); q.pop()) {
    for (int x = q.front(); int u : g[x])
      if (s[u] == -1) {
         if (pre[u] = x, s[u] = 1, match[u] == n) {
           for (int a = u, b = x, last;
    b!= n; a = last, b = pre[a])
```

q.push(match[u]); s[match[u]] = 0; } else if (!s[u] && Find(u) != Find(x)) {

int 1 = LCA(u, x);
Blossom(x, u, 1); Blossom(u, x, 1);

 $\label{eq:matching} \texttt{Matching}(\underbrace{\textbf{int}}_{-n}) \; : \; \mathsf{n}(_\mathsf{n}), \; \mathsf{fa}(\mathsf{n}+1), \; \mathsf{s}(\mathsf{n}+1), \; \mathsf{vis}(\mathsf{n}$

+ 1), pre(n + 1, n), match(n + 1, n), g(n) {}

last = match[b], match[b] = a, match[a] =

```
3.10 Flow Model
```

return false;

int solve() { int ans = 0;

return ans;

void add_edge(int u, int v)

{ g[u].pb(v), g[v].pb(u); }

for (int x = 0; x < n; ++x)

if (match[x] == n) ans += Bfs(x);

} // match[x] == n means not matched

return true;

Maximum/Minimum flow with lower bound / Circulation problem

- 1. Construct super source S and sink T.
- 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l.
- 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
- 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
- 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
- 1. Redirect every edge: $y \to x$ if $(x,y) \in M$, $x \to y$ otherwise.
- 2. DFS from unmatched vertices in X.
- 3. $x \in X$ is chosen iff x is unvisited.
- 4. $y \in Y$ is chosen iff y is visited.
- · Minimum cost cyclic flow
- 1. Consruct super source ${\cal S}$ and sink ${\cal T}$
- 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c > 0, otherwise connect $y \to x$ with (cost, cap) = (-c, 1)
- 3. For each edge with c<0 , sum these cost as K , then increase d(y) by 1, decrease d(x) by 1
- 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0,d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) = (0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
- 1. Binary search on answer, suppose we're checking answer ${\it T}$
- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source $s \to v$, $v \in G$ with capacity K
- 4. For each edge (u,v,w) in G , connect $u\to v$ and $v\to u$ with capacity w
- 5. For $v \in G$, connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w + 2w(v))$
- 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
 - 1. Let $w'(u,v)=w(u,v)-\mu(u)-\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
 - 2. Find the minimum weight matching M with w'. The answer is $\sum \mu(v) + w'(M).$
- Project selection problem
 - 1. If $p_v>0$, create edge (s,v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$.
- 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- Dual of minimum cost maximum flow
 - 1. Capacity c_{uv} , Flow f_{uv} , Cost w_{uv} , Required Flow difference for vertex b_u .
- 2. If all w_{uv} are integers, then optimal solution can happen when all p_u are integers.

$$\min \sum_{uv} w_{uv} f_{uv}$$

$$-f_{uv} \ge -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv})$$

$$\sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u}$$

$$p_{u} \ge 0$$

4 Geometry

4.1 Geometry Template [86f0f1]

```
using ld = ll;
using pdd = pair<ld, ld>;
#define X first
#define Y second
// Ld eps = 1e-7;

pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator-(pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y}; }
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
```

```
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(ld v){ return v > eps ? 1 : ( v < -eps ? -1
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity(p, a, b) && sgn(dot(a - p, b - p)
    ) <= 0; }
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  if(btw(p1, p3, p4) || btw(p2, p3, p4) || btw(p3, p1,
       p2) || btw(p4, p1, p2))
    return true;
  return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&
    ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
  ld a124 = cross(p2 - p1, p4 - p1);
return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(
    p2 - p1); }
(pdd \text{ reflection(pdd p1, pdd p2, pdd p3)})
  return p3 + perp(p2 - p1) * cross(p3 - p1, p2 - p1) /
      abs2(p2 - p1) * 2; }
pdd linearTransformation(pdd p0, pdd p1, pdd q0, pdd q1
     , pdd r) {
  pdd dp = p1 - p0, dq = q1 - q0, num(cross(dp, dq),
       dot(dp, dq));
  return q0 + pdd(cross(r - p0, num), dot(r - p0, num))
        / abs2(dp);
} // from line p0--p1 to q0--q1, apply to r
4.2 Polar Angle Comparator [808e89]
```

4.3 Minkowski Sum [b3028c]

```
return ans;
```

4.4 Intersection of Circle and Convex Polygon [63653d]

```
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
   S = (C/2)*r*r;
   h = a*b*sin(C)/c;
    if (h < r \&\& B < PI/2) S -= (acos(h/r)*r*r - h*sqrt
        (r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
   S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
 return S;
double areaPolyCircle(const vector<pdd> poly,const pdd
    &O.const double r){
  double S=0;
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=\_area(poly[i]-0,poly[(i+1)\%SZ(poly)]-0,r)*ori(0,f)
        poly[i],poly[(i+1)%SZ(poly)]);
 return fabs(S);
}
```

4.5 Intersection of Circles [f7a2fe]

4.6 Tangent Line of Circles [c51d90]

```
vector<Line> CCtang( const Cir& c1 , const Cir& c2 ,
    int sign1 ){
  vector<Line> ret;
  double d_sq = abs2( c1.0 - c2.0 );
  if (sgn(d_sq) == 0) return ret;
  double d = sqrt(d_sq);
  pdd v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d; // cos t
  if (c * c > 1) return ret;
  double h = sqrt(max( 0.0, 1.0 - c * c)); // sin t
  for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
        v.Y * c + sign2 * h * v.X);
    pdd p1 = c1.0 + n * c1.R;
    pdd p2 = c2.0 + n * (c2.R * sign1);
    if (sgn(p1.X - p2.X) == 0 and
        sgn(p1.Y - p2.Y) == 0)
      p2 = p1 + perp(c2.0 - c1.0);
    ret.pb(Line(p1, p2));
  }
  return ret:
}
```

4.7 Intersection of Line and Convex Polygon [157258]

```
int TangentDir(vector<pll> &C, pll dir) {
  return cyc_tsearch(SZ(C), [&](int a, int b) {
    return cross(dir, C[a]) > cross(dir, C[b]);
```

```
});
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pll a, pll b, vector<pll> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 || cmpL(B) > 0)
    return pii(-1, -1); // no collision
  auto gao = [&](int 1, int r) {
    for (int t = 1; (1 + 1) % n != r; ) {
      int m = ((1 + r + (1 < r? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(t) ? 1 : r) = m;
    return (1 + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
  if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!cmpL(res.X) && !cmpL(res.Y)) // along side i, i
    switch ((res.X - res.Y + n + 1) % n) {
      case 0: return pii(res.X, res.X);
      case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned in the same order as the line hits the
      convex */
  return res;
} // convex cut: (r, l]
```

4.8 Intersection of Line and Circle [9183db]

```
vector<pdd> circleLineIntersection(pdd c, double r, pdd
    a, pdd b) {
  pdd p = a + (b - a) * dot(c - a, b - a) / abs2(b - a)
  ;
  double s = cross(b - a, c - a), h2 = r * r - s * s /
    abs2(b - a);
  if (sgn(h2) < 0) return {};
  if (sgn(h2) == 0) return {p};
  pdd h = (b - a) / abs(b - a) * sqrt(h2);
  return {p - h, p + h};
}</pre>
```

4.9 Point in Circle [ecf954]

```
// return q's relation with circumcircle of tri(p[0],p
       [1],p[2])
bool in_cc(const array<pll, 3> &p, pll q) {
    __int128 det = 0;
    for (int i = 0; i < 3; ++i)
       det += __int128(abs2(p[i]) - abs2(q)) * cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
    return det > 0; // in: >0, on: =0, out: <0
}</pre>
```

4.10 Point in Convex [82b81e]

4.11 Half Plane Intersection [d34e39]

```
auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128) a02Y * a12X - (__int128) a02X *
      a12Y > 0;
/* Having solution, check size > 2 */
/* --^-- Line.X --^-- Line.Y --^-- */
vector<Line> halfPlaneInter(vector<Line> arr) {
 sort(iter(arr), [&](Line a, Line b) -> int {
  if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
    return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  auto pop_back = [&](int t, Line p) {
    while (SZ(dq) >= t \& !isin(p, dq[SZ(dq) - 2], dq.
         back()))
      dq.pop_back();
  auto pop_front = [&](int t, Line p) {
    while (SZ(dq) >= t \&\& !isin(p, dq[0], dq[1]))
      dq.pop_front();
  for (auto p : arr)
    if (cmp(dq.back().Y - dq.back().X, p.Y - p.X, 0) !=
          -1)
      pop_back(2, p), pop_front(2, p), dq.pb(p);
  pop_back(3, dq[0]), pop_front(3, dq.back());
  return vector<Line>(iter(dq));
```

4.12 HPI General Line [043534]

```
using i128 = __int128;
struct LN {
  11 a, b, c; // ax + by + c <= 0
  pll dir() const { return pll(a, b); }
  LN(11 ta, 11 tb, 11 tc) : a(ta), b(tb), c(tc) {}
  LN(pll S, pll T): a((T-S).Y), b(-(T-S).X), c(cross(T, T-S).X)
      S)) {}
pdd intersect(LN A, LN B) {
  double c = cross(A.dir(), B.dir());
  i128 a = i128(A.c) * B.a - i128(B.c) * A.a;
i128 b = i128(A.c) * B.b - i128(B.c) * A.b;
  return pdd(-b / c, a / c);
bool cov(LN 1, LN A, LN B) {
  i128 c = cross(A.dir(), B.dir());
  i128 \ a = i128(A.c) * B.a - i128(B.c) * A.a;
  i128 b = i128(A.c) * B.b - i128(B.c) * A.b;
  return sign(a * 1.b - b * 1.a + c * 1.c) * sign(c) >=
       0;
bool operator<(LN a, LN b) {</pre>
  if (int c = cmp(a.dir(), b.dir(), false); c != -1)
      return c;
  return i128(abs(b.a) + abs(b.b)) * a.c > i128(abs(a.a
      ) + abs(a.b)) * b.c;
}
```

4.13 Minimum Enclosing Circle [5af6d5]

```
return {c, r};
4.14 3D Point [badbbd]
struct Point {
  double x, y, z;
  Point(double _x = 0, double _y = 0, double _z = 0): x = (_x), y(_y), z(_z){}
  Point(pdd p) { x = p.X, y = p.Y, z = abs2(p); }
Point operator-(Point p1, Point p2)
{ return Point(p1.x - p2.x, p1.y - p2.y, p1.z - p2.z);
Point operator+(Point p1, Point p2)
{ return Point(p1.x + p2.x, p1.y + p2.y, p1.z + p2.z);
Point operator*(Point p1, double v)
{ return Point(p1.x * v, p1.y * v, p1.z * v); }
Point operator/(Point p1, double v)
{ return Point(p1.x / v, p1.y / v, p1.z / v); }
Point cross(Point p1, Point p2)
{ return Point(p1.y * p2.z - p1.z * p2.y, p1.z * p2.x -
     p1.x * p2.z, p1.x * p2.y - p1.y * p2.x); }
double dot(Point p1, Point p2)
{ return p1.x * p2.x + p1.y * p2.y + p1.z * p2.z; }
double abs(Point a)
{ return sqrt(dot(a, a)); }
Point cross3(Point a, Point b, Point c)
{ return cross(b - a, c - a); }
double area(Point a, Point b, Point c)
{ return abs(cross3(a, b, c)); }
double volume(Point a, Point b, Point c, Point d)
{ return dot(cross3(a, b, c), d - a); }
//Azimuthal angle (longitude) to x-axis in interval [-
    pi, pi]
double phi(Point p) { return atan2(p.y, p.x); }
//Zenith angle (latitude) to the z-axis in interval [0,
     pi1
double theta(Point p) { return atan2(sqrt(p.x * p.x + p
    .y * p.y), p.z); }
Point masscenter(Point a, Point b, Point c, Point d)
\{ return (a + b + c + d) / 4; \}
pdd proj(Point a, Point b, Point c, Point u) {
// proj. u to the plane of a, b, and c
  Point e1 = b - a;
  Point e2 = c - a;
  e1 = e1 / abs(e1);
  e2 = e2 - e1 * dot(e2, e1);
  e2 = e2 / abs(e2);
  Point p = u - a;
  return pdd(dot(p, e1), dot(p, e2));
Point rotate_around(Point p, double angle, Point axis)
  double s = sin(angle), c = cos(angle);
  Point u = axis / abs(axis);
  return u * dot(u, p) * (1 - c) + p * c + cross(u, p)
}
4.15 ConvexHull3D [156311]
struct convex hull 3D {
struct Face {
  int a, b, c;
  Face(int ta, int tb, int tc): a(ta), b(tb), c(tc) {}
}; // return the faces with pt indexes
vector<Face> res;
vector<Point> P:
convex_hull_3D(const vector<Point> &_P): res(), P(_P) {
// all points coplanar case will WA, O(n^2)
  int n = SZ(P);
  if (n <= 2) return; // be careful about edge case</pre>
  // ensure first 4 points are not coplanar
  swap(P[1],\ *find\_if(iter(P),\ [\&](auto\ p)\ \{\ return\ sgn
      (abs2(P[0] - p)) != 0; }));
  swap(P[2],\ *find\_if(iter(P),\ [\&](auto\ p)\ \{\ return\ sgn
      (abs2(cross3(p, P[0], P[1]))) != 0; }));
  swap(P[3], *find_if(iter(P), [&](auto p) { return sgn
```

(volume(P[0], P[1], P[2], p)) != 0; }));
vector<vector<int>> flag(n, vector<int>(n));

res.emplace_back(0, 1, 2); res.emplace_back(2, 1, 0);

```
for (int i = 3; i < n; ++i) {
                                                                 divide(0, n - 1);
    vector<Face> next;
    for (auto f : res) {
      int d = sgn(volume(P[f.a], P[f.b], P[f.c], P[i]))
      if (d <= 0) next.pb(f);</pre>
      int ff = (d > 0) - (d < 0);
      flag[f.a][f.b] = flag[f.b][f.c] = flag[f.c][f.a]
    for (auto f : res) {
      auto F = [\&](int x, int y) {
        if (flag[x][y] > 0 \&\& flag[y][x] \leftarrow 0)
          next.emplace_back(x, y, i);
      F(f.a, f.b); F(f.b, f.c); F(f.c, f.a);
    res = next:
 }
bool same(Face s, Face t) {
                                                                   }
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.a])) != 0)
                                                                   return false;
       return 0;
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.b])) != 0)
  if (sgn(volume(P[s.a], P[s.b], P[s.c], P[t.c])) != 0)
                                                                 while (true) {
       return 0;
  return 1;
int polygon_face_num() {
  int ans = 0;
  for (int i = 0; i < SZ(res); ++i)</pre>
    ans += none_of(res.begin(), res.begin() + i, [&](
                                                                            id])))
        Face g) { return same(res[i], g); });
  return ans;
double get_volume() {
  double ans = 0;
  for (auto f : res)
    ans += volume(Point(0, 0, 0), P[f.a], P[f.b], P[f.c
        ]);
  return fabs(ans / 6);
                                                                     else ++it;
double get_dis(Point p, Face f) {
  Point p1 = P[f.a], p2 = P[f.b], p3 = P[f.c];
                                                               }
  double a = (p2.y - p1.y) * (p3.z - p1.z) - (p2.z - p1
                                                            };
      .z) * (p3.y - p1.y);
  double b = (p2.z - p1.z) * (p3.x - p1.x) - (p2.x - p1
      .x) * (p3.z - p1.z);
  double c = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1
      .y) * (p3.x - p1.x);
  double d = 0 - (a * p1.x + b * p1.y + c * p1.z);
  return fabs(a * p.x + b * p.y + c * p.z + d) / sqrt(a
       * a + b * b + c * c);
}
};
// n^2 delaunay: facets with negative z normal of
// convexhull of (x, y, x^2 + y^2), use a pseudo-point // (0, 0, \inf) to avoid degenerate case
                                                                        - pts[u]);
4.16 Delaunay Triangulation [6a9916]
/* Delaunay Triangulation:
                                                            }
   Given a sets of points on 2D plane, find a
   triangulation such that no points will strictly
   inside circumcircle of any triangle. */
struct Edge {
  int id; // oidx[id]
  list<Edge>::iterator twin;
  Edge(int _id = 0):id(_id) {}
                                                               ld res = 0;
struct Delaunay { // 0-base
 int n;
  vector<int> oidx;
  vector<list<Edge>> head; // result udir. graph
  vector<pll> p;
  Delaunay(int _n, vector<pll> _p): n(_n), oidx(n),
      head(n), p(n) {
    iota(iter(oidx), 0);
    for (int i = 0; i < n; ++i) head[i].clear();</pre>
    sort(iter(oidx), [&](int a, int b)
        { return _p[a] < _p[b]; });
```

for (int i = 0; i < n; ++i) p[i] = _p[oidx[i]];</pre>

```
void addEdge(int u, int v) {
    head[u].push_front(Edge(v));
    head[v].push_front(Edge(u));
    head[u].begin()->twin = head[v].begin();
    head[v].begin()->twin = head[u].begin();
  void divide(int 1, int r) {
    if (1 == r) return;
    if (l + 1 == r) return addEdge(l, l + 1);
    int mid = (1 + r) \gg 1, nw[2] = \{1, r\};
    divide(l, mid), divide(mid + 1, r);
    auto gao = [&](int t)
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      for (auto it : head[nw[t]]) {
        int v = ori(pt[1], pt[0], p[it.id]);
        if (v > 0 || (v == 0 && abs2(pt[t ^ 1] - p[it.
            id]) < abs2(pt[1] - pt[0]))) \\
          return nw[t] = it.id, true;
    while (gao(0) || gao(1));
    addEdge(nw[0], nw[1]); // add tangent
      pll pt[2] = {p[nw[0]], p[nw[1]]};
      int ch = -1, sd = 0;
      for (int t = 0; t < 2; ++t)
        for (auto it : head[nw[t]])
          if (ori(pt[0], pt[1], p[it.id]) > 0 && (ch ==
               -1 || in_cc({pt[0], pt[1], p[ch]}, p[it.
            ch = it.id, sd = t;
      if (ch == -1) break; // upper common tangent
      for (auto it = head[nw[sd]].begin(); it != head[
          nw[sd]].end(); )
        if (seg_strict_intersect(pt[sd], p[it->id], pt[
            sd ^ 1], p[ch]))
          head[it->id].erase(it->twin), head[nw[sd]].
              erase(it++);
      nw[sd] = ch, addEdge(nw[0], nw[1]);
4.17 Voronoi Diagram [e4f408]
// all coord. is even, you may want to call
    halfPlaneInter after then
vector<vector<Line>> vec;
void build_voronoi_line(int n, vector<pll> &pts) {
  Delaunay tool(n, pts); // Delaunay
  vec.clear(), vec.resize(n);
  for (int i = 0; i < n; ++i)</pre>
    for (auto e : tool.head[i]) {
      int u = tool.oidx[i], v = tool.oidx[e.id];
      pll m = (pts[v] + pts[u]) / 2LL, d = perp(pts[v])
      vec[u].pb(Line(m, m + d));
4.18 Polygon Union [9fbf66]
ld rat(pll a, pll b) {
  return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
 // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
      pll A = p[a], B = p[(a + 1) \% SZ(p)];
      vector<pair<ld, int>> segs = {{0, 0}, {1, 0}};
      for (auto &q : poly) {
        if (&p == &q) continue;
        for (int b = 0; b < SZ(q); ++b) {
  pll C = q[b], D = q[(b + 1) % SZ(q)];</pre>
          int sc = ori(A, B, C), sd = ori(A, B, D);
          if (sc != sd && min(sc, sd) < 0) {</pre>
            1d sa = cross(D - C, A - C), sb = cross(D -
                 C, B - C);
```

```
segs.pb(sa / (sa - sb), sgn(sc - sd));
          if (!sc && !sd && &q < &p && sgn(dot(B - A, D</pre>
            - C)) > 0) {
segs.pb(rat(C - A, B - A), 1);
            segs.pb(rat(D - A, B - A), -1);
          }
        }
      sort(iter(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      1d sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {</pre>
        if (!cnt) sum += segs[j].X - segs[j - 1].X;
        cnt += segs[j].Y;
      res += cross(A, B) * sum;
    }
  return res / 2;
4.19 Tangent Point to Convex Hull [523bc1]
/* The point should be strictly out of hull
```

```
return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
  auto gao = [&](int s) {
   return cyc_tsearch(SZ(C), [&](int x, int y)
   { return ori(p, C[x], C[y]) == s; });
 return pii(gao(1), gao(-1));
} // return (a, b), ori(p, C[a], C[b]) >= 0
```

4.20 Heart [082d19]

```
pdd circenter(pdd p0, pdd p1, pdd p2) { // 156d1f
 p1 = p1 - p0, p2 = p2 - p0; // radius = abs(center)
  double x1 = p1.X, y1 = p1.Y, x2 = p2.X, y2 = p2.Y;
  double m = 2. * (x1 * y2 - y1 * x2);
 pdd center;
  center.X = (x1 * x1 * y2 - x2 * x2 * y1 + y1 * y2 * (
     y1 - y2)) / m;
 center.Y = (x1 * x2 * (x2 - x1) - y1 * y1 * x2 + x1 *
      y2 * y2) / m;
  return center + p0;
pdd incenter(pdd p1, pdd p2, pdd p3) { // radius = area
     / s * 2
  double a = abs(p2 - p3), b = abs(p1 - p3), c = abs(p1
       - p2);
  double s = a + b + c;
  return (a * p1 + b * p2 + c * p3) / s;
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter(p1, p2, p3) * 3 - circenter(p1, p2,
```

4.21 Rotating Sweep Line [f5f689]

```
struct Event {
  pll d; int u, v;
  bool operator<(const Event &b) const {</pre>
    int ret = cmp(d, b.d, false);
    return ret == -1 ? false : ret; } // no tie-break
void rotatingSweepLine(const vector<pll> &p) {
  const int n = SZ(p);
  vector<Event> e; e.reserve(n * (n - 1));
  for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++) // pos[i] < pos[j] when
         the event occurs
      if (i != j) e.pb(p[j] - p[i], i, j);
  sort(iter(e));
  vector<int> ord(n), pos(n);
  iota(iter(ord), 0);
  sort(iter(ord), [&](int i, int j) { // initial order
      return p[i].Y != p[j].Y ? p[i].Y < p[j].Y : p[i].</pre>
          X < p[j].X; \});
  for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
  // initialize
```

```
11
  for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
    // do something
    vector<pii> tmp;
    for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
      tmp.pb(pii(e[j].u, e[j].v));
    sort(iter(tmp), [&](pii x, pii y){
        return pii(pos[x.ff], pos[x.ss]) < pii(pos[y.ff</pre>
            ], pos[y.ss]); });
    for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
      tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
        make_tuple(ord[pos[y]], ord[pos[x]], pos[y],
            pos[x]);
  }
}
4.22 Vector In Poly [c6d0fa]
// ori(a, b, c) >= 0, valid: "strict" angle from a-b to
  return ori(a, b, p) >= strict && ori(a, p, c) >=
      strict:
```

```
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
// whether vector{cur, p} in counter-clockwise order
    prv, cur, nxt
bool inside(pll prv, pll cur, pll nxt, pll p, int
   strict) {
  if (ori(cur, nxt, prv) >= 0)
    return btwangle(cur, nxt, prv, p, strict);
  return !btwangle(cur, prv, nxt, p, !strict);
```

4.23 Convex Hull DP [6dc001]

```
sort(iter(pts), [&](pll x, pll y) {
  return x.Y != y.Y ? x.Y < y.Y : x.X < y.X;
auto getvec = [&](pii x) { return pts[x.ss] - pts[x.ff
    ]; };
vector<pii> trans;
for (int j = 0; j < n; j++)</pre>
  for (int k = 0; k < n; k++)
    if (j != k) trans.pb(pii(j, k));
sort(iter(trans), [&](pii x, pii y) -> bool{
  int tmp = cmp(getvec(x), getvec(y), false);
  if (tmp != -1) return tmp;
  pll v = getvec(x);
  return dot(v, pts[x.ff]) > dot(v, pts[y.ff]);
// DP for convex hull vertices (no points on edges)
auto solve = [\&](int bottom) { // <math>O(n^3)
  // vector<ll> dp(n);
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether bottom -> j is legal
    // init trans -> j
  for (auto [i, j] : trans) {
    if (i <= bottom || j <= bottom ||</pre>
        ori(pts[bottom], pts[i], pts[j]) <= 0) continue</pre>
    // check whether i -> j is legal
    // normal trans i -> j
  for (int j = bottom + 1; j < n; j++) {</pre>
    // check whether j -> bottom is legal
    // end trans j ->
  }
};
for(int i = 0; i < n; i++) solve(i);</pre>
```

4.24 Calculate Points in Triangle [bf746f]

```
// all points are distinct
// cnt[i][j] = # of point k s.t. strictly above ij, and
     i < k < j
// cnt2[i][j] = # of points k s.t. strictly in ij
// preprocess space: O(n^2), time: O(n^3), query time:
    0(1)
vector cnt(n, vector<int>(n)), cnt2(n, vector<int>(n));
for (int i = 0; i < n; i++)</pre>
  for (int j = 0; j < n; j++){
    if (pts[i] >= pts[j]) continue;
    for (int k = 0; k < n; k++) {
```

```
if (pts[i] < pts[k] && pts[k] < pts[j]) {</pre>
        int tmp = ori(pts[i], pts[j], pts[k]);
        if (tmp > 0) cnt[i][j]++; // only for i < j
        else if (tmp == 0) cnt2[i][j]++, cnt2[j][i]++;
   }
 }
auto calc_tri = [&](array<int, 3> arr) { // strictly
  sort(iter(arr), [&](int x, int y){ return pts[x] <</pre>
      pts[y]; });
  auto [x, y, z] = arr;
  int tmp = ori(pts[x], pts[y], pts[z]);
  if (tmp == 0) return 0;
 else if (tmp < 0)</pre>
    return cnt[x][z] - cnt[x][y] - cnt[y][z] - cnt2[x][
        y] - cnt2[y][z] - 1;
  else return cnt[x][y] + cnt[y][z] - cnt[x][z] - cnt2[
      x][z];
};
```

Graph 5

5.1 BCC [d04ebe]

```
struct BCC{ // O-based, allow multi edges but not allow
  int n, m, cnt = 0;
  // n:|V|, m:|E|, cnt:#bcc
  // bcc i : vertices bcc_v[i] and edges bcc_e[i]
  vector<vector<int>> bcc_v, bcc_e;
  vector<vector<pii>>> g; // original graph
  vector<pii> edges; // 0-based
  BCC(int _n, vector<pii> _edges):
    n(_n), m(SZ(_edges)), g(_n), edges(_edges){
      for(int i = 0; i < m; i++){</pre>
        auto [u, v] = edges[i];
        g[u].pb(pii(v, i)); g[v].pb(pii(u, i));
      }
  void make_bcc(){ bcc_v.pb(); bcc_e.pb(); cnt++; }
  // modify these if you need more information
void add_v(int v){ bcc_v.back().pb(v); }
  void add_e(int e){ bcc_e.back().pb(e); }
  void build(){
    vector\langle int \rangle in(n, -1), low(n, -1), stk;
    vector<vector<int>> up(n);
    int ts = 0;
    auto _dfs = [&](auto dfs, int now, int par, int pe)
          -> void{
      if(pe != -1) up[now].pb(pe);
      in[now] = low[now] = ts++;
      stk.pb(now);
      for(auto [v, e] : g[now]){
        if(e == pe) continue;
        if(in[v] != -1){
          if(in[v] < in[now]) up[now].pb(e);</pre>
          low[now] = min(low[now], in[v]);
          continue;
        dfs(dfs, v, now, e);
        low[now] = min(low[now], low[v]);
      if((now != par && low[now] >= in[par]) || (now ==
           par && SZ(g[now]) == 0)){
        make_bcc();
        for(int v = stk.back();; v = stk.back()){
          stk.pop_back(), add_v(v);
          for(int e : up[v]) add_e(e);
          if(v == now) break;
        if(now != par) add_v(par);
      }
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) _dfs(_dfs, i, i, -1);
 }
};
```

5.2 SCC [2c9a01]

```
struct SCC{ // 0-based, output reversed topo order
 int n, cnt = 0;
```

```
12
  vector<vector<int>> g;
  vector<int> sccid;
  explicit SCC(int _n): n(_n), g(n), sccid(n, -1) {}
  void add_edge(int u, int v){
    g[u].pb(v);
  void build(){
    vector<int> in(n, -1), low(n), stk;
    vector<bool> instk(n);
    int ts = 0;
    auto dfs1 = [&](auto dfs, int now) -> void{
      stk.pb(now); instk[now] = true;
      in[now] = low[now] = ts++;
      for(int i : g[now]){
        if(in[i] == -1)
          dfs(dfs, i), low[now] = min(low[now], low[i])
        else if(instk[i] && in[i] < in[now])</pre>
          low[now] = min(low[now], in[i]);
      if(low[now] == in[now]){
        for(; stk.back() != now; stk.pop_back())
          sccid[stk.back()] = cnt, instk[stk.back()] =
              false;
        sccid[now] = cnt++, instk[now] = false, stk.
            pop_back();
      }
    };
    for(int i = 0; i < n; i++)</pre>
      if(in[i] == -1) dfs1(dfs1, i);
  }
};
5.3 2-SAT [0686a5]
struct SAT { // 0-based
  int n;
  vector<bool> istrue;
  SCC scc;
  SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
  int neg(int a) {
    return a >= n ? a - n : a + n;
  void add_clause(int a, int b) {
    scc.add_edge(neg(a), b), scc.add_edge(neg(b), a);
  bool solve() {
    scc.build();
    for (int i = 0; i < n; ++i) {</pre>
      if (scc.sccid[i] == scc.sccid[i + n]) return
      istrue[i] = scc.sccid[i] < scc.sccid[i + n];</pre>
      istrue[i + n] = !istrue[i];
    return true;
  }
};
5.4
      Dominator Tree [2da9bb]
struct Dominator {
  vector<vector<int>> g, r, rdom; int tk;
  vector<int> dfn, rev, fa, sdom, dom, val, rp;
  Dominator(int_n): n(n), g(n), r(n), rdom(n), tk(0)
    dfn = rev = fa = sdom = dom =
      val = rp = vector<int>(n, -1); }
  void add_edge(int x, int y) { g[x].push_back(y); }
```

```
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk; tk++;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
  if (int p = find(fa[x], 1); p != -1) {
    if (sdom[val[x]] > sdom[val[fa[x]]])
      val[x] = val[fa[x]];
```

```
return x==a[x] ? x : a[x] = Y(Y, a[x]); };
      fa[x] = p;
      return c ? p : val[x];
                                                               auto S = [&](int i) { return o(o, e[i].s); };
    } else return c ? fa[x] : val[x];
                                                               int pc = v[root] = n;
                                                               for (int i = 0; i < n; ++i) if (v[i] == -1)</pre>
  vector<int> build(int s) {
                                                                  for (int p = i; v[p]<0 || v[p]==i; p = S(r[p])) {</pre>
    // return the father of each node in dominator tree
                                                                    if (v[p] == i)
    dfs(s); // p[i] = -2 if i is unreachable, par[s] =
                                                                      for (int q = pc++; p != q; p = S(r[p])) {
        -1
                                                                        h[p].tag -= h[p].top().v; h[q].join(h[p]);
    for (int i = tk - 1; i >= 0; --i) {
                                                                        pa[p] = a[p] = q;
      for (int u : r[i])
        sdom[i] = min(sdom[i], sdom[find(u)]);
                                                                    while (S(h[p].top().i) == p) h[p].pq.pop();
                                                                    v[p] = i; r[p] = h[p].top().i;
      if (i) rdom[sdom[i]].push_back(i);
      for (int u : rdom[i]) {
                                                                 }
        int p = find(u);
                                                               vector<int> ans;
        dom[u] = (sdom[p] == i ? i : p);
                                                               for (int i = pc - 1; i >= 0; i--) if (v[i] != n) {
                                                                  for (int f = e[r[i]].t; f!=-1 && v[f]!=n; f = pa[f
      if (i) merge(i, rp[i]);
                                                                    v[f] = n;
    vector < int > p(n, -2); p[s] = -1;
                                                                 ans.push_back(r[i]);
    for (int i = 1; i < tk; ++i)</pre>
      if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
                                                               return ans; // default minimize, returns edgeid array
    for (int i = 1; i < tk; ++i)</pre>
                                                             }
      p[rev[i]] = rev[dom[i]];
                                                             5.7 Vizing [58a6ca]
    return p;
                                                             // find D+1 edge coloring of a graph with max deg D, O(
}:
                                                                  nm)
                                                             struct Vizing { // returns maxdeg+1 edge coloring in
5.5 Virtual Tree [6abeb5]
                                                                  adjacent matrix G
                                                               int n; // 1-based for vertices and colors, simple
vector<int> vG[N];
                                                                    graph
int top, st[N];
                                                               vector<vector<int>> C. G:
int vrt = -1;
                                                               vector<int> X, vst;
void insert(int u) {
                                                               Vizing(int _n): n(_n),
  if (top == -1) return st[++top] = vrt = u, void();
                                                               C(n + 1, vector < int > (n + 2)), G(n + 1, vector < int > (n + 1))
  int p = LCA(st[top], u);
                                                                    + 1)),
    if(dep[vrt] > dep[p]) vrt = p;
                                                               X(n + 1, 1), vst(n + 1) {}
  if (p == st[top]) return st[++top] = u, void();
                                                               void solve(vector<pii> &E) {
  while (top >= 1 && dep[st[top - 1]] >= dep[p])
                                                                  auto update = [&](int u)
    vG[st[top - 1]].pb(st[top]), --top;
                                                                 { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
auto color = [&](int u, int v, int c) {
  if (st[top] != p)
    vG[p].pb(st[top]), --top, st[++top] = p;
                                                                    int p = G[u][v];
  st[++top] = u;
                                                                    G[u][v] = G[v][u] = c;
                                                                    C[u][c] = v, C[v][c] = u;
void reset(int u) {
                                                                    C[u][p] = C[v][p] = 0;
  for (int i : vG[u]) reset(i);
                                                                    if (p) X[u] = X[v] = p;
  vG[u].clear();
                                                                    else update(u), update(v);
                                                                   return p:
void solve(vector<int> &v) {
                                                                  };
  top = -1;
                                                                  auto flip = [&](int u, int c1, int c2) {
  sort(iter(v),
                                                                    int p = C[u][c1];
      [&](int a, int b) { return dfn[a] < dfn[b]; });</pre>
                                                                    swap(C[u][c1], C[u][c2]);
  for (int i : v) insert(i);
                                                                    if (p) G[u][p] = G[p][u] = c2;
  while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
                                                                    if (!C[u][c1]) X[u] = c1;
  // do something
                                                                    if (!C[u][c2]) X[u] = c2;
  reset(vrt);
                                                                    return p;
                                                                  for (int t = 0; t < SZ(E); ++t) {</pre>
5.6 Fast DMST [7b274d]
                                                                    int u = E[t].ff, v0 = E[t].ss, v = v0, c0 = X[u],
struct E { int s, t; ll w; }; // 0-base
                                                                         c = c0, d;
struct PQ {
                                                                    vector<pii> L;
  struct P {
                                                                    fill(iter(vst), 0);
    11 v; int i;
                                                                    while (!G[u][v0]) {
                                                                      L.emplace_back(v, d = X[v]);
    bool operator>(const P &b) const { return v > b.v;
                                                                      if (!C[v][c]) for (int a = SZ(L) - 1; a >= 0;
                                                                           --a) c = color(u, L[a].ff, c);
                                                                      else if (!C[u][d]) for (int a = SZ(L) - 1; a >=
  priority_queue<P, vector<P>, greater<>> pq; 11 tag;
      // min heap
                                                                           0; --a) color(u, L[a].ff, L[a].ss);
                                                                      else if (vst[d]) break;
  void push(P p) { p.v -= tag; pq.emplace(p); }
  P top() { P p = pq.top(); p.v += tag; return p; }
                                                                      else vst[d] = 1, v = C[u][d];
  void join(PQ &b) {
    if (pq.size() < b.pq.size())</pre>
                                                                    if (!G[u][v0]) {
                                                                      for (; v; v = flip(v, c, d), swap(c, d));
      swap(pq, b.pq), swap(tag, b.tag);
    while (!b.pq.empty()) push(b.top()), b.pq.pop();
                                                                      if (int a; C[u][c0]) {
                                                                        for (a = SZ(L) - 2; a >= 0 && L[a].ss != c;
}; // O(E log^2 V), use leftist tree for O(E log V)
                                                                             --a);
vector<int> dmst(const vector<E> &e, int n, int root) {
                                                                        for (; a >= 0; --a) color(u, L[a].ff, L[a].ss
  vector<PQ> h(n * 2);
for (int i = 0; i < int(e.size()); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
                                                                      else --t;
  vector<int> a(n * 2); iota(iter(a), 0);
vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
                                                                    }
                                                                 }
  auto o = [\&](auto Y, int x) \rightarrow int {
```

5.8 Maximum Clique [1ad4b2]

};

```
struct MaxClique { // fast when N <= 100
 bitset<N> G[N], cs[N];
  int ans, sol[N], q, cur[N], d[N], n;
  void init(int _n) {
   n = n;
    for (int i = 0; i < n; ++i) G[i].reset();</pre>
  void add_edge(int u, int v) {
   G[u][v] = G[v][u] = 1;
  void pre_dfs(vector<int> &r, int 1, bitset<N> mask) {
    if (1 < 4) {
      for (int i : r) d[i] = (G[i] & mask).count();
      sort(iter(r), [\&](int x, int y) \{ return d[x] > d \}
          [y]; });
    }
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
     int k = 1;
      while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
      cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
      for (int p = cs[k]._Find_first(); p < N; p = cs[k</pre>
          ]._Find_next(p))
        r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<int> &r, vector<int> &c, int 1,
      bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
      r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
      for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, 1, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
   }
  int solve() {
   vector<int> r(n);
    ans = q = 0, iota(iter(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans:
};
```

5.9 Number of Maximal Clique [11fa26]

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
  if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
int tsn = 0, tnn = 0;
      copy_n(all[d], an, all[d + 1]);
       all[d + 1][an] = v;
       for (int j = 0; j < sn; ++j)</pre>
         if (g[v][some[d][j]])
```

```
some[d + 1][tsn++] = some[d][j];
for (int j = 0; j < nn; ++j)
    if (g[v][none[d][j]])
        none[d + 1][tnn++] = none[d][j];
    dfs(d + 1, an + 1, tsn, tnn);
    some[d][i] = 0, none[d][nn++] = v;
}
int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
}
};</pre>
```

5.10 Minimum Mean Cycle [3e5d2b]

```
11 road[N][N]; // input here
struct MinimumMeanCycle {
  11 dp[N + 5][N], n;
  pll solve() {
    ll a = -1, b = -1, L = n + 1;
    for (int i = 2; i <= L; ++i)</pre>
       for (int k = 0; k < n; ++k)
         for (int j = 0; j < n; ++j)
           dp[i][j] =
              min(dp[i - 1][k] + road[k][j], dp[i][j]);
    for (int i = 0; i < n; ++i) {</pre>
       if (dp[L][i] >= INF) continue;
       11 ta = 0, tb = 1;
for (int j = 1; j < n; ++j)</pre>
         if (dp[j][i] < INF &&</pre>
           ta * (L - j) < (dp[L][i] - dp[j][i]) * tb)
ta = dp[L][i] - dp[j][i], tb = L - j;
       if (ta == 0) continue;
       if (a == -1 || a * tb > ta * b) a = ta, b = tb;
    if (a != -1) {
      11 g = 
                __gcd(a, b);
       return pll(a / g, b / g);
    return pll(-1LL, -1LL);
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)</pre>
       for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;
};
```

5.11 Minimum Steiner Tree [21aceal

```
// O(V 3^T + V^2 2^T)
struct SteinerTree { // 0-base
  static const int T = 10, N = 105, INF = 1e9;
  int n, dst[N][N], dp[1 << T][N], tdst[N];
int vcost[N]; // the cost of vertexs</pre>
  void init(int _n) {
    n = _n;
    for (int i = 0; i < n; ++i) {</pre>
      for (int j = 0; j < n; ++j) dst[i][j] = INF;</pre>
      dst[i][i] = vcost[i] = 0;
  void add_edge(int ui, int vi, int wi) {
    dst[ui][vi] = min(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)
      for (int i = 0; i < n; ++i)
         for (int j = 0; j < n; ++j)</pre>
           dst[i][j] =
             min(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int> &ter) {
    shortest_path();
    int t = SZ(ter);
    for (int i = 0; i < (1 << t); ++i)</pre>
      for (int j = 0; j < n; ++j) dp[i][j] = INF;</pre>
    for (int i = 0; i < n; ++i) dp[0][i] = vcost[i];</pre>
    for (int msk = 1; msk < (1 << t); ++msk) {</pre>
      if (!(msk & (msk - 1))) {
        int who = __lg(msk);
```

```
National Taiwan University fruit advantages
        for (int i = 0; i < n; ++i)</pre>
          dp[msk][i] =
            vcost[ter[who]] + dst[ter[who]][i];
      for (int i = 0; i < n; ++i)</pre>
        for (int submsk = (msk - 1) & msk; submsk;
             submsk = (submsk - 1) \& msk)
          dp[msk][i] = min(dp[msk][i],
            dp[submsk][i] + dp[msk ^ submsk][i] -
              vcost[i]);
      for (int i = 0; i < n; ++i) {</pre>
        tdst[i] = INF;
        for (int j = 0; j < n; ++j)</pre>
          tdst[i] =
            min(tdst[i], dp[msk][j] + dst[j][i]);
      for (int i = 0; i < n; ++i) dp[msk][i] = tdst[i];</pre>
    int ans = INF;
    for (int i = 0; i < n; ++i)</pre>
     ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
 }
};
5.12 Count Cycles [c7e8f2]
// ord = sort by deg decreasing, rk[ord[i]] = i
// D[i] = edge point from rk small to rk big
for (int x : ord) { // c3
 for (int y : D[x]) vis[y] = 1;
  for (int y : D[x]) for (int z : D[y]) c3 += vis[z];
  for (int y : D[x]) vis[y] = 0;
for (int x : ord) { // c4
 for (int y : D[x]) for (int z : adj[y])
   if (rk[z] > rk[x]) c4 += vis[z]++;
  for (int y : D[x]) for (int z : adj[y])
    if (rk[z] > rk[x]) --vis[z];
} // both are O(M*sqrt(M))
    Math
      Extended Euclidean Algorithm [c51ae9]
// ax+ny = 1, ax+ny == ax == 1 \ (mod \ n)
void extgcd(ll x, ll y, ll &g, ll &a, ll &b) {
 if (y == 0) g = x, a = 1, b = 0;
  else extgcd(y, x % y, g, b, a), b -= (x / y) * a;
6.2 Floor & Ceil [134881]
ll ifloor(ll a, ll b){
 return a / b - (a % b && (a < 0) ^ (b < 0));
il iceil(ll a, ll b){
 return a / b + (a % b && (a < 0) ^ (b > 0));
6.3 Legendre [4e4b23]
// the Jacobi symbol is a generalization of the
    Legendre symbol,
// such that the bottom doesn't need to be prime.
// (n|p) -> same as legendre
// (n|ab) = (n|a)(n|b)
// work with long long
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
```

if (a == 0) **return** 0;

if (a & m & 2) s = -s;

// -1: a isn't a quad res of p

a >>= r;

return s;

// 0: a == 0

swap(a, m);

const int r = __builtin_ctz(a);

if ((r & 1) && ((m + 2) & 4)) s = -s;

```
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  if(int jc = Jacobi(a, p); jc <= 0) return jc;</pre>
  int b, d;
  for (;;) {
    b = rand() \% p;
     d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
  int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
  for (int e = (1LL + p) >> 1; e; e >>= 1) {
    if (e & 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 %
             p)) % p;
       g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
       g0 = tmp;
     tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)
    )) % p;
f1 = (2LL * f0 * f1) % p;
     f0 = tmp;
  }
  return g0;
}
6.4 Simplex [aa7741]
// maximize c^T x
// subject to Ax <= b, x >= 0
// and stores the solution;
typedef long double T; // long double, Rational, double
      + mod<P>...
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-9, inf = 1/.0;
#define ltj(X) if(s == -1 || mp(X[j],N[j]) < mp(X[s],N[
    s1)) s=i
#define rep(i, 1, n) for(int i = 1; i < n; i++)
struct LPSolver {
  int m, n;
  vector<int> N, B;
  vvd D;
   \begin{array}{l} LPSolver(\textbf{const} \ vvd\& \ A, \ \textbf{const} \ vd\& \ b, \ \textbf{const} \ vd\& \ c) \ : \\ m(SZ(b)), \ n(SZ(c)), \ N(n+1), \ B(m), \ D(m+2, \ vd(n+2)) \ \{ \end{array} 
       rep(i,0,m) rep(j,0,n) D[i][j] = A[i][j];
       rep(i,0,m) { B[i] = n+i; D[i][n] = -1; D[i][n+1]
            = b[i];}
        \begin{split} & \text{rep}(j,0,n) ~ \{ ~ N[j] = j; ~ D[m][j] = -c[j]; ~ \} \\ & N[n] = -1; ~ D[m+1][n] = 1; \end{split} 
  void pivot(int r, int s) {
     T *a = D[r].data(), inv = 1 / a[s];
     rep(i,0,m+2) if (i != r \&\& abs(D[i][s]) > eps) {
       T *b = D[i].data(), inv2 = b[s] * inv;
       rep(j,0,n+2) b[j] -= a[j] * inv2;
       b[s] = a[s] * inv2;
     rep(j,0,n+2) if (j != s) D[r][j] *= inv;
     rep(i,0,m+2) if (i != r) D[i][s] *= -inv;
    D[r][s] = inv;
     swap(B[r], N[s]);
  bool simplex(int phase) {
     int x = m + phase - 1;
     for (;;) {
       int s = -1:
       rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
       if (D[x][s] >= -eps) return true;
       int r = -1;
       rep(i,0,m) {
         if (D[i][s] <= eps) continue;
if (r == -1 || mp(D[i][n+1] / D[i][s], B[i])</pre>
              < mp(D[r][n+1] / D[r][s], B[r])) r = i;
       if (r == -1) return false;
       pivot(r, s);
```

// else: return X with $X^2 \% p == a$

// doesn't work with long long

bool prime(ll n){ // 8859aa

```
vector<ll> tmp = {2, 325, 9375, 28178, 450775,
                                                                        9780504, 1795265022};
                                                                    for(ll i : tmp)
  T solve(vd &x) {
                                                                      if(!Miller_Rabin(i, n)) return false;
    int r = 0;
                                                                    return true:
    rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] < -eps) {</pre>
                                                                  map<ll, int> cnt;
       pivot(r, n);
                                                                  void PollardRho(ll n) { // 173531
                                                                    if (n == 1) return;
       if (!simplex(2) || D[m+1][n+1] < -eps) return -</pre>
                                                                    if (prime(n)) return ++cnt[n], void();
           inf:
       rep(i,0,m) if (B[i] == -1) {
                                                                    if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2],
         int s = 0;
         rep(j,1,n+1) ltj(D[i]);
                                                                    11 x = 2, y = 2, d = 1, p = 1;
         pivot(i, s);
                                                                  #define f(x, n, p) ((mul(x, x, n) + p) % n)
       }
                                                                    while (true) {
                                                                      if (d != n && d != 1) {
                                                                        PollardRho(n / d);
    bool ok = simplex(1); x = vd(n);
    rep(i,0,m) if (B[i] < n) \times [B[i]] = D[i][n+1];
                                                                        PollardRho(d);
    return ok ? D[m][n+1] : inf;
                                                                        return;
                                                                      if (d == n) ++p;
};
                                                                      x = f(x, n, p), y = f(f(y, n, p), n, p);
      Simplex Construction
                                                                      d = gcd(abs(x - y), n);
Standard form: maximize \sum_{1 \leq i \leq n} c_i x_i such that \sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j for
all 1 \le j \le m and x_i \ge 0 for all 1 \le i \le n.
1. In case of minimization, let c'_i = -c_i
                                                                  6.8 XOR Basis [006505]
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
3. \sum_{1 \leq i \leq n}^{-} A_{ji} x_i = b_j \rightarrow \mathsf{add} \subseteq \mathsf{and} \supseteq.
                                                                  const int digit = 60; // [0, 2^digit)
4. If x_i has no lower bound, replace x_i with x_i - x_i'
                                                                  struct Basis{
                                                                    int total = 0, rank = 0;
6.6 DiscreteLog [da27bf]
                                                                    vector<11> b;
                                                                    Basis(): b(digit) {}
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
                                                                    bool add(ll v){ // Gauss Jordan Elimination
  unordered_map<int, int> p;
  int b = 1;
                                                                      for(int i = digit - 1; i >= 0; i--){
  for (int i = 0; i < kStep; ++i) {</pre>
                                                                        if(!(1LL << i & v)) continue;</pre>
                                                                        if(b[i] != 0){
    p[y] = i;
    y = 1LL * y * x % m;
                                                                           v ^= b[i];
    b = 1LL * b * x % m;
                                                                           continue;
                                                                        for(int j = 0; j < i; j++)</pre>
  for (int i = 0; i < m + 10; i += kStep) {</pre>
    s = 1LL * s * b % m;
                                                                          if(1LL << j & v) v ^= b[j];</pre>
                                                                         for(int j = i + 1; j < digit; j++)</pre>
    if (p.find(s) != p.end()) return i + kStep - p[s];
                                                                          if(1LL << i & b[j]) b[j] ^= v;</pre>
                                                                        b[i] = v;
  return -1;
                                                                        rank++;
int DiscreteLog(int x, int y, int m) {
                                                                        return true;
  if (m == 1) return 0;
  int s = 1;
                                                                      return false;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
s = 1LL * s * x % m;
                                                                    11 \text{ getmax}(11 \text{ x} = 0){
                                                                      for(ll i : b) x = max(x, x ^ i);
                                                                      return x;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
                                                                    11 \text{ getmin}(11 \times = 0){
                                                                      for(ll i : b) x = min(x, x ^ i);
  if (fpow(x, p, m) != y) return -1;
  return p; //returns: x^p = y \pmod{m}
                                                                      return x;
                                                                    bool can(ll x){
6.7 Miller Rabin & Pollard Rho [d3ecd2]
                                                                      return getmin(x) == 0;
// n < 4,759,123,141
                             3: 2, 7, 61
                                                                    11 kth(11 k){ // kth smallest, 0-indexed
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
                                                                      vector<ll> tmp;
// n < 3,474,749,660,383 6 : primes <= 13
                                                                      for(ll i : b) if(i) tmp.pb(i);
// n < 2^64
                                                                      11 \text{ ans} = 0;
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
                                                                      for(int i = 0; i < SZ(tmp); i++)</pre>
11 mul(l1 a, l1 b, l1 n){
                                                                        if(1LL << i & k) ans ^= tmp[i];</pre>
  return (__int128)a * b % n;
                                                                      return ans;
                                                                   }
bool Miller_Rabin(ll a, ll n) { // 06308c
                                                                 };
  if ((a = a % n) == 0) return 1;
  if (n % 2 == 0) return n == 2;
                                                                  6.9 Linear Equation [056191]
  11 \text{ tmp} = (n - 1) / ((n - 1) & (1 - n));
                                                                  vector<int> RREF(vector<vector<11>> &mat) { // 9cd26b
  11 t = _{lg(((n - 1) \& (1 - n))), x = 1;}
  for (; tmp; tmp >>= 1, a = mul(a, a, n))
                                                                    int N = SZ(mat), M = SZ(mat[0]);
    if (tmp & 1) x = mul(x, a, n);
                                                                    int rk = 0;
  if (x == 1 || x == n - 1) return 1;
                                                                    vector<int> cols;
  while (--t)
                                                                    for (int i = 0; i < M; i++) {</pre>
    if ((x = mul(x, x, n)) == n - 1) return 1;
                                                                      int cnt = -1;
                                                                      for (int j = N - 1; j >= rk; j--)
  return 0;
```

if(mat[j][i] != 0) cnt = j;

if (cnt == -1) continue;

```
swap(mat[rk], mat[cnt]);
    11 lead = mat[rk][i];
    for (int j = 0; j < M; j++) mat[rk][j] = mat[rk][j]</pre>
         * inv(lead) % MOD;
    for (int j = 0; j < N; j++) {
      if (j == rk) continue;
      11 tmp = mat[j][i];
      for (int k = 0; k < M; k++)</pre>
        mat[j][k] = (mat[j][k] - mat[rk][k] * tmp % MOD
              + MOD) % MOD;
    cols.pb(i);
   rk++;
  return cols;
// sol = particualr + linear combination of homogenous
struct LinearEquation { // 2702e2
  bool ok;
  vector<11> par; //particular solution (Ax = b)
  vector<vector<ll>> homo; //homogenous (Ax = \theta)
  vector<vector<ll>> rref;
  //first M columns are matrix A
  //last column of eq is vector b
  void solve(const vector<vector<ll>> &eq) {
    int M = SZ(eq[0]) - 1;
    rref = eq;
    auto piv = RREF(rref);
    int rk = piv.size();
    if(piv.size() && piv.back() == M)
      return ok = 0, void();
    ok = 1:
    par.resize(M);
    vector<bool> ispiv(M);
    for (int i = 0;i < rk;i++) {</pre>
      par[piv[i]] = rref[i][M];
      ispiv[piv[i]] = 1;
    for (int i = 0; i < M; i++) {</pre>
      if (ispiv[i]) continue;
      vector<ll> h(M);
      h[i] = 1;
      for (int j = 0; j < rk; j++)</pre>
        h[piv[j]] = rref[j][i] ? MOD - rref[j][i] : 0;
      homo.pb(h);
 }
};
```

6.10 Chinese Remainder Theorem [6ef4a3]

6.11 Sqrt Decomposition [8d7bc0]

```
// for all i in [l, r], floor(n / i) = x
for(int l = 1, r; l <= n; l = r + 1){
   int x = ifloor(n, l);
   r = ifloor(n, x);
}
// for all i in [l, r], ceil(n / i) = x
for(int l, r = n; r >= 1; r = l - 1){
   int x = iceil(n, r);
   l = iceil(n, x);
}
```

6.12 Floor Sum

```
• m = \lfloor \frac{an+b}{c} \rfloor
• Time complexity: O(\log n)
```

```
f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
                                       \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)}{2} + \left\lfloor \frac{b}{c} \right\rfloor \cdot (n+1)
                                      +f(a \bmod c, b \bmod c, c, n),
                                                                                                  a \ge c \lor b \ge c
                                                                                                  n < 0 \lor a = 0
                                     nm - f(c, c - b - 1, a, m - 1), otherwise
                            \left( \left\lfloor \frac{a}{c} \right\rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor \cdot \frac{n(n+1)}{2} \right)
                            +g(a \bmod c, b \bmod c, c, n),
                                                                                                            a \ge c \lor b \ge c
                                                                                                            n<0\vee a=0
                              \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1))
                            -h(c, c-b-1, a, m-1),
                                                                                                            otherwise
h(a,b,c,n) = \sum_{i=1}^{n} \lfloor \frac{ai+b}{i} \rfloor^{2}
                            \left\{ \left\lfloor \frac{a}{c} \right\rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \left\lfloor \frac{b}{c} \right\rfloor^2 \cdot (n+1) \right\}
                             +\lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1)
                             +h(a \bmod c, b \bmod c, c, n)
                            +2\lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n)
                             +2\lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n),
                                                                                                           a \ge c \lor b \ge c
                                                                                                            n < 0 \lor a = 0
                             nm(m+1) - 2g(c, c-b-1, a, m-1)
                             -2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), otherwise
```

7 Polynomial

7.1 FWHT [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)
invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)
    for (int i = 0; i < n; i += L)</pre>
      for (int j = i; j < i + (L >> 1); ++j)
         a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[N][1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N</pre>
    ];
void subset_convolution(int *a, int *b, int *c, int L)
  // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;</pre>
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)</pre>
  fwt(f[i], n, 1), fwt(g[i], n, 1);
for (int i = 0; i <= L; ++i)</pre>
     for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
h[i][x] += f[j][x] * g[i - j][x];</pre>
  for (int i = 0; i <= L; ++i) fwt(h[i], n, -1);</pre>
  for (int i = 0; i < n; ++i) c[i] = h[ct[i]][i];</pre>
```

7.2 FFT [13ec2f]

```
// Errichto: FFT for double works when the result < 1
    e15, and < 1e18 with long double

using val_t = complex<double>;
template<int MAXN>
struct FFT {
    const double PI = acos(-1);
    val_t w[MAXN];
    FFT() {
        for (int i = 0; i < MAXN; ++i) {
            double arg = 2 * PI * i / MAXN;
            w[i] = val_t(cos(arg), sin(arg));
        }
    }
}</pre>
```

#define V (*this)

template <int MAXN, 11 P, 11 RT> // MAXN = 2^k

struct Poly : vector<ll> { // coefficients in [0, P)

```
void bitrev(vector<val_t> &a, int n) //same as NTT
                                                                using vector<11>::vector;
  void trans(vector<val_t> &a, int n, bool inv = false)
                                                                static inline NTT<MAXN, P, RT> ntt;
                                                                int n() const { return (int)size(); } // n() >= 1
    bitrev(a, n);
for (int L = 2; L <= n; L <<= 1) {</pre>
                                                                Poly(const Poly &p, int m) : vector<ll>(m) { copy_n(p
                                                                    .data(), min(p.n(), m), data()); }
      int dx = MAXN / L, dl = L >> 1;
                                                                Poly &irev() { return reverse(data(), data() + n()),
      for (int i = 0; i < n; i += L) {</pre>
                                                                    V; }
                                                                Poly &isz(int m) { return resize(m), V; }
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
                                                                static ll minv(ll x) { return ntt.minv(x); }
          val_t = a[j + dl] * (inv ? conj(w[x]) : w
                                                              // == fb1867 ==
                                                                Poly &iadd(const Poly &rhs) { // db5668
               [x]);
          a[j + dl] = a[j] - tmp;
                                                                  fi(0, n()) if ((V[i] += rhs[i]) >= P) V[i] -= P;
                                                                  return V; // need n() == rhs.n()
          a[j] += tmp;
      }
                                                                Poly &imul(11 k) { // a8df26
                                                                  fi(0, n()) V[i] = V[i] * k % P;
    if (inv) {
                                                                  return V;
      for (int i = 0; i < n; ++i) a[i] /= n;</pre>
    }
                                                                Poly Mul(const Poly &rhs) const { // 46caf3
                                                                  int m = 1;
  //multiplying two polynomials A * B:
                                                                  while (m < n() + rhs.n() - 1) m <<= 1;</pre>
  //fft.trans(A, siz, 0), fft.trans(B, siz, 0):
                                                                  assert(m <= MAXN);</pre>
  //A[i] *= B[i], fft.trans(A, siz, 1);
                                                                  Poly X(V, m), Y(rhs, m);
};
                                                                  ntt(X, m), ntt(Y, m);
                                                                  fi(0, m) X[i] = X[i] * Y[i] % P;
7.3
      NTT [bf683f]
                                                                  ntt(X, m, true);
                                                                  return X.isz(n() + rhs.n() - 1);
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
                                                                Poly Inv() const { // 796a37
//1255*(2^20)+1, 1315962881, 3
                                                                  if (n() == 1) return {minv(V[0])};
//51*(2^25)+1, 1711276033, 29
                                                                  int m = 1; // need V[0] != 0, 2*sz<=MAXN</pre>
// only works when sz(A) + sz(B) - 1 <= MAXN
                                                                  while (m < n() * 2) m <<= 1;</pre>
template<int MAXN, 11 P, 11 RT> //MAXN must be 2^k
                                                                  assert(m <= MAXN);</pre>
struct NTT {
                                                                  Poly Xi = Poly(V, (n() + 1) / 2).Inv().isz(m);
  11 w[MAXN];
                                                                  Poly Y(V, m);
  11 mpow(ll a, ll n);
                                                                  ntt(Xi, m), ntt(Y, m);
  11 minv(ll a) { return mpow(a, P - 2); }
                                                                  fi(0, m) {
  NTT() {
                                                                    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    ll dw = mpow(RT, (P - 1) / MAXN);
                                                                    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
    w[0] = 1;
    for (int i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw
                                                                  ntt(Xi, m, true);
         % P;
                                                                  return Xi.isz(n());
  void bitrev(vector<ll> &a, int n) {
                                                                Poly &shift_inplace(const 11 &c) { // 0c04f6
    int i = 0;
                                                                  int n = V.n(); // 2 * sz <= MAXN
vector<ll> fc(n), ifc(n);
    for (int j = 1; j < n - 1; ++j) {
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
                                                                  fc[0] = ifc[0] = 1;
      if (j < i) swap(a[i], a[j]);</pre>
                                                                  for (int i = 1; i < n; i++) {
  fc[i] = fc[i - 1] * i % P;</pre>
    }
  }
                                                                    ifc[i] = minv(fc[i]);
  void operator()(vector<ll> &a, int n, bool inv =
      false) { //0 <= a[i] < P
                                                                  for (int i = 0; i < n; i++) V[i] = V[i] * fc[i] % P</pre>
    bitrev(a, n);
    for (int L = 2; L <= n; L <<= 1) {</pre>
                                                                  Poly g(n);
      int dx = MAXN / L, dl = L >> 1;
                                                                  11 cp = 1;
      for (int i = 0; i < n; i += L) {</pre>
                                                                  for (int i = 0; i < n; i++) g[i] = cp * ifc[i] % P,</pre>
        for (int j = i, x = 0; j < i + d1; ++j, x += dx
                                                                       cp = cp * c % P;
                                                                  V = V.irev().Mul(g).isz(n).irev();
          ll tmp = a[j + dl] * w[x] % P;
                                                                  for (int i = 0; i < n; i++) V[i] = V[i] * ifc[i] %</pre>
          if ((a[j + d1] = a[j] - tmp) < 0) a[j + d1]
                                                                  return V:
          if ((a[j] += tmp) >= P) a[j] -= P;
                                                               }
        }
                                                             // == 7b2835 ==
      }
                                                                Poly shift(const ll &c) const { return Poly(V).
                                                                    shift_inplace(c); }
    if (inv) {
                                                                Poly _Sqrt() const { // Jacobi(V[0], P) = 1
      reverse(a.begin()+1, a.begin()+n);
                                                                  if (n() == 1) return {QuadraticResidue(V[0], P)};
      11 invn = minv(n);
                                                                  Poly X = Poly(V, (n() + 1) / 2)._Sqrt().isz(n());
      for (int i = 0; i < n; ++i) a[i] = a[i] * invn %</pre>
                                                                  return X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 +
                                                                      1);
                                                               }
  }
                                                             // == b46641 ==
};
                                                                Poly Sqrt() const { // 1aa942
                                                                  Poly a; // 2 * sz <= MAXN
      Polynomial Operation [77a8a8]
                                                                  bool has = 0;
                                                                  for (int i = 0; i < n(); i++) {</pre>
// == b4233a ==
#define fi(s, n) for (int i = (int)(s); i < (int)(n);
                                                                    if (V[i]) has = 1;
                                                                    if (has) a.push_back(V[i]);
    ++i)
#define neg(x) (x ? P - x : 0)
```

if (!has) return V;

return Poly();

if ((n() + a.n()) % 2 || Jacobi(a[0], P) != 1) {

```
a = a.isz((n() + a.n()) / 2)._Sqrt();
    int sz = a.n();
    a.isz(n());
    rotate(a.begin(), a.begin() + sz, a.end());
  pair<Poly, Poly> DivMod(const Poly &rhs) const { // 5
    if (n() < rhs.n()) return {{0}, V};
const int m = n() - rhs.n() + 1;</pre>
    Poly X(rhs); // (rhs.)back() != 0
    X.irev().isz(m);
    Poly Y(V);
    Y.irev().isz(m);
    Poly Q = Y.Mul(X.Inv()).isz(m).irev();
   X = rhs.Mul(Q), Y = V;
fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
    return {Q, Y.isz(max(1, rhs.n() - 1))};
// == 76b1af ==
 Poly Dx() const {
    Poly ret(n() - 1);
    fi(0, ret.n()) ret[i] = (i + 1) * V[i + 1] % P;
    return ret.isz(max(1, ret.n()));
  Poly Sx() const {
    Poly ret(n() + 1);
    fi(0, n()) ret[i + 1] = minv(i + 1) * V[i] % P;
    return ret;
  Poly _tmul(int nn, const Poly &rhs) const {
    Poly Y = Mul(rhs).isz(n() + nn - 1);
    return Poly(Y.data() + n() - 1, Y.data() + Y.n());
// == 3afa3f ==
  vector<ll> _eval(const vector<ll> &x, const vector<</pre>
      Poly> &up) const { // fb6553
    const int m = (int)x.size();
    if (!m) return {};
    vector<Poly> down(m * 2);
    // down[1] = DivMod(up[1]).second;
    // fi(2, m * 2) down[i] = down[i / 2].DivMod(up[i])
        .second;
    down[1] = Poly(up[1]).irev().isz(n()).Inv().irev().
        _tmul(m, V);
    fi(2, m * 2) down[i] = up[i ^ 1]._tmul(up[i].n() -
        1, down[i / 2]);
    vector<11> y(m);
    fi(0, m) y[i] = down[m + i][0];
    return y;
  static vector<Poly> _tree1(const vector<ll> &x) { //
      f5c433
    const int m = (int)x.size();
    vector<Poly> up(m * 2);
    fi(0, m) up[m + i] = {neg(x[i]), 1};
    for (int i = m - 1; i > 0; --i) up[i] = up[i * 2].
        Mul(up[i * 2 + 1]);
    return up;
  }
  vector<ll> Eval(const vector<ll> &x) const { // 1e5,
    auto up = _tree1(x);
    return _eval(x, up);
  static Poly Interpolate(const vector<11> &x, const
      vector<ll> &y) { // d7bae4
    const int m = (int)x.size(); // 1e5, 1.4s
    vector<Poly> up = _{tree1(x), down(m * 2)};
    vector<ll> z = up[1].Dx()._eval(x, up);
    fi(0, m) z[i] = y[i] * minv(z[i]) % P;
    fi(0, m) down[m + i] = {z[i]};
    for (int i = m - 1; i > 0; --i)
  down[i] = down[i * 2].Mul(up[i * 2 + 1]).iadd(
          down[i * 2 + 1].Mul(up[i * 2]));
    return down[1]:
// == c066ab ==
 Poly Ln() const { //V[0] == 1, 2*sz <= MAXN
    return Dx().Mul(Inv()).Sx().isz(n());
```

```
Poly Exp() const { //V[0] == 0,2*sz <= MAXN
    if (n() == 1) return {1};
    Poly X = Poly(V, (n() + 1) / 2).Exp().isz(n());
    Poly Y = X.Ln();
    Y[0] = P - 1;
    fi(0, n()) if ((Y[i] = V[i] - Y[i]) < 0) Y[i] += P;
    return X.Mul(Y).isz(n());
// == 3f1d86 ==
  //M := P(P - 1). If k >= M, k := k % M + M.
  Poly Pow(11 k) const { // 2*sz<=MAXN // d08261
    int nz = 0;
    while (nz < n() && !V[nz]) ++nz;</pre>
    if (nz * min(k, (11)n()) >= n()) return Poly(n());
    if (!k) return Poly(Poly{1}, n());
    Poly X(data() + nz, data() + nz + n() - nz * k);
const ll c = ntt.mpow(X[0], k % (P - 1));
    return X.Ln().imul(k % P).Exp().imul(c).irev().isz(
        n()).irev();
  // sum_j w_j [x^j] f(x^i) for i \in [0, m]
  Poly power_projection(Poly wt, int m) { // 277119
    assert(n() == wt.n()); // 4*sz <= MAXN!
    if (!n()) {
      return Poly(m + 1, 0);
    if (V[0] != 0) {
      11 c = V[0];
      V[0] = 0;
      Poly A = V.power_projection(wt, m);
      fi(0, m + 1) A[i] = A[i] * fac[i] % P; //
          factorial
      Poly B(m + 1);
      11 pow = 1;
      fi(0, m + 1) B[i] = pow * ifac[i] % P, pow = pow
           * c % P; // inv. of fac
      A = A.Mul(B).isz(m + 1);
      fi(0, m + 1) A[i] = A[i] * fac[i] % P;
      return A;
    int n = 1;
    while (n < V.n()) n *= 2;</pre>
    isz(n), wt.isz(n).irev();
    int k = 1;
    Poly p(wt, 2 * n), q(V, 2 * n);
    q.imul(P - 1);
    while (n > 1) {
      Poly r(2 * n * k);
      fi(0, 2 * n * k) r[i] = (i % 2 == 0 ? q[i] : neg(
           q[i]));
      Poly pq = p.Mul(r).isz(4 * n * k);
      Poly qq = q.Mul(r).isz(4 * n * k);
      fi(0, 2 * n * k) {
        pq[2 * n * k + i] += p[i];
        qq[2 * n * k + i] += q[i] + r[i];
        pq[2 * n * k + i] %= P;
        qq[2 * n * k + i] %= P;
      fill(p.begin(), p.end(), 0);
      fill(q.begin(), q.end(), 0);
for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {</pre>
        p[n * j + i] = pq[(2 * n) * j + (2 * i + 1)];

q[n * j + i] = qq[(2 * n) * j + (2 * i + 0)];
      n /= 2, k *= 2;
    Poly ans(k);
    fi(0, k) ans[i] = p[2 * i];
    return ans.irev().isz(m + 1);
  Poly FPSinv() { // 2c54b4
    const int n = V.n() - 1;
    if (n == -1) return {};
    assert(V[0] == 0);
    if (n == 0) return V;
    assert(V[1] != 0);
    ll c = V[1], ic = minv(c);
    imul(ic);
    Poly wt(n + 1);
    wt[n] = 1;
```

```
Poly A = V.power_projection(wt, n);
    Poly g(n);
    fi(1, n + 1) g[n - i] = n * A[i] % P * minv(i) % P;
    g = g.Pow(neg(minv(n)));
    g.insert(g.begin(), 0);
    11 pow = 1;
    return g;
  Poly TMul(const Poly &rhs) const { // this[i] - rhs[j
       ] = k; // 7b552c
    return Poly(*this).irev().Mul(rhs).isz(n()).irev();
  Poly FPScomp(Poly g) { // solves V(g(x)) // 332bb2
    auto rec = [&](auto &rec, int n, int k, Poly Q) ->
         Poly {
       if (n == 1) {
         Poly p(2 * k);
         irev();
         fi(0, k) p[2 * i] = V[i];
         return p;
      Poly R(2 * n * k);
      fi(0, 2 * n * k) R[i] = (i % 2 == 0 ? Q[i] : neg(
           Q[i]));
      Poly QQ = Q.Mul(R).isz(4 * n * k);
fi(0, 2 * n * k) {
 QQ[2 * n * k + i] += Q[i] + R[i];
         QQ[2 * n * k + i] \% = P;
      Poly nxt_Q(2 * n * k);
      for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
         nxt_Q[n * j + i] = QQ[(2 * n) * j + (2 * i + 0)
      Poly nxt_p = rec(rec, n / 2, k * 2, nxt_Q);
      Poly pq(4 * n * k);
      for(int j = 0; j < 2 * k; j++) fi(0, n / 2) {
   pq[(2 * n) * j + (2 * i + 1)] += nxt_p[n * j +
             i];
         pq[(2 * n) * j + (2 * i + 1)] \% = P;
      Poly p(2 * n * k);
      fi(0, 2 * n * k) p[i] = (p[i] + pq[2 * n * k + i]

 % P;

      pq.pop_back();
      Poly x = pq.TMul(R);
      fi(0, 2 * n * k) p[i] = (p[i] + x[i]) % P;
    };
    int sz = 1;
    while(sz < n() || sz < g.n()) sz <<= 1;</pre>
    return isz(sz), rec(rec, sz, 1, g.imul(P-1).isz(2 *
          sz)).isz(sz).irev();
 }
};
#undef fi
#undef V
#undef neg
using Poly_t = Poly<1 << 19, 998244353, 3>;
Ordinary Generating Function
• C(x) = A(rx): c_n = r^n a_n 的一般生成函數。
• C(x) = A(x) + B(x): c_n = a_n + b_n 的一般生成函數。
• C(x) = A(x)B(x): c_n = \sum_{i=0}^n a_i b_{n-i} 的一般生成函數。
• C(x) = A(x)^k: c_n = \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k} 的一般生成函數。
```

7.5 Generating Function

- C(x) = xA(x)': $c_n = na_n$ 的一般生成函數。
- $C(x) = \frac{A(x)}{1-x}$: $c_n = \sum_{i=0}^n a_i$ 的一般生成函數。
- $C(x) = A(1) + x \frac{A(1) A(x)}{1 x}$: $c_n = \sum_{i=n}^{\infty} a_i$ 的一般生成函數。

常用展開式

- $\frac{1}{1-x} = 1 + x + x^2 + \ldots + x^n + \ldots$
- $(1+x)^a = \sum_{n=0}^{\infty} {n \choose n} x^n$, ${n \choose n} = \frac{a(a-1)(a-2)...(a-n+1)}{n!}$.

常見生函

卡特蘭數: $f(x) = \frac{1 - \sqrt{1 - 4x}}{2}$

Exponential Generating Function

 a_0, a_1, \ldots 的指數生成函數:

$$\hat{A}(x) = \sum_{i=0}^{\infty} \frac{a_i}{i!} = a_0 + a_1 x + \frac{a_2}{2!} x^2 + \frac{a_3}{3!} x^3 + \dots$$

- $\hat{C}(x) = \hat{A}(x) + \hat{B}(x)$: $c_n = a_n + b_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}^{(k)}(x)$: $c_n = a_{n+k}$ 的指數生成函數
- $\hat{C}(x) = x\hat{A}(x)$: $c_n = na_n$ 的指數生成函數
- $\hat{C}(x) = \hat{A}(x)$. $\hat{C}_n = han \text{ 5.13 MeV}$ 知识 的指數生成函數 $\hat{C}(x) = \hat{A}(x)\hat{B}(x)$: $c_n = \sum_{k=0}^n \binom{n}{i} a_k b_{n-k}$ 的指數生成函數 $\hat{C}(x) = \hat{A}(x)^k$: $\sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_i a_{i_2} \dots a_{i_k}$ 的指數生成函數
- $\hat{C}(x) = \exp(A(x))$: 假設 A(x) 是一個分量 (component) 的生成函數,那 $\hat{C}(x)$ 是將 n 個有編號的東西分成若干個分量的指數生成函數

Lagrange's Inversion Formula

如果 F 跟 G 互反,則有 F(0),G(0)=0, $F'(0),G'(0)\neq 0$ 。若 H 為任意 FPS,則

$$n[x^n]G(x) = [x^{n-1}] \frac{1}{(F(x)/x)^n}$$
$$n[x^n]H(G(x)) = [x^{n-1}]H'(x) \frac{1}{(F(x)/x)^n}$$

7.6 Bostan Mori [41c3bc]

const 11 mod = 998244353;

```
NTT<262144, mod, 3> ntt;
// Finds the k-th coefficient of P / Q in O(d log d log
// size of NTT has to > 2 * d
11 BostanMori(vector<11> P, vector<11> Q, long long k)
  int d = max((int)P.size(), (int)Q.size() - 1);
  vector M = \{P, Q\};
  M[0].resize(d, 0);
  M[1].resize(d + 1, 0);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  vector<ll> Qn(sz);
  vector N(2, vector<ll>(sz));
  while(k) {
    fill(iter(Qn), 0);
    for(int i = 0; i < d + 1; i++){
  Qn[i] = M[1][i] * ((i & 1) ? -1 : 1);</pre>
      if(Qn[i] < 0) Qn[i] += mod;</pre>
    ntt(Qn, sz, false);
    11 t[2] = \{k \& 1, 0\};
    for(int i = 0; i < 2; i++){</pre>
      fill(iter(N[i]), 0);
      copy(iter(M[i]), N[i].begin());
      ntt(N[i], sz, false);
for(int j = 0; j < sz; j++)</pre>
        N[i][j] = N[i][j] * Qn[j] % mod;
      ntt(N[i], sz, true);
      for(int j = t[i]; j < 2 * siz(M[i]); j += 2){</pre>
        M[i][j >> 1] = N[i][j];
    k \gg 1;
  return M[0][0] * ntt.minv(M[1][0]) % mod;
11 LinearRecursion(vector<ll> a, vector<ll> c, ll k) {
    // a_n = \sum_{j=1}^{d} c_j a_{n-j}
  int d = siz(a);
  int sz = (2 * d + 1 == 1 ? 2 : (1 << (__lg(2 * d) +
      1)));
  c[0] = mod - 1;
  for(l1 &i : c) i = i ? mod - i : 0;
  auto A = a; A.resize(sz);
  auto C = c; C.resize(sz);
  ntt(A, sz, false), ntt(C, sz, false);
  for(int i = 0; i < sz; i++) A[i] = A[i] * C[i] % mod;</pre>
  ntt(A, sz, true);
  A.resize(d);
```

```
return BostanMori(A, c, k);
8
    String
8.1 KMP Algorithm [c8b75f]
// 0-based
// fail[i] = max k < i s.t. s[0..k] = s[i-k..i]
vector<int> kmp_build_fail(const string &s){
 int n = SZ(s);
  vector<int> fail(n, -1);
  int cur = -1:
  for(int i = 1; i < n; i++){</pre>
    while(cur != -1 && s[cur + 1] != s[i])
      cur = fail[cur];
    if(s[cur + 1] == s[i])
      cur++;
    fail[i] = cur;
 return fail:
void kmp_match(const string &s, const vector<int> &fail
    , const string &t){
  int cur = -1;
 int n = SZ(s), m = SZ(t);
  for(int i = 0; i < m; i++){</pre>
    while(cur != -1 && (cur + 1 == n || s[cur + 1] != t
      cur = fail[cur];
    if(cur + 1 < n \&\& s[cur + 1] == t[i])
      cur++;
    // cur = max \ k \ s.t. \ s[0..k] = t[i-k..i]
}
/* center i: radius z[i * 2 + 1] / 2
   both aba, abba have radius 2 */
  string s = "%";
  int 1 = 0, r = 0;
```

8.2 Manacher Algorithm [caf0f4]

```
center i, i + 1: radius z[i * 2 + 2] / 2
vector<int> manacher(const string &tmp){ // 0-based
  for(char c : tmp) s += c, s += '%';
  vector<int> z(SZ(s));
  for(int i = 0; i < SZ(s); i++){</pre>
   z[i] = r > i ? min(z[2 * 1 - i], r - i) : 1;
    while(i - z[i] >= 0 \&\& i + z[i] < SZ(s)
           && s[i + z[i]] == s[i - z[i]])
      ++z[i];
    if(z[i] + i > r) r = z[i] + i, l = i;
 }
  return z;
```

8.3 Lyndon Factorization [7c612b]

```
// partition s = w[0] + w[1] + ... + w[k-1],
// w[0] >= w[1] >= ... >= w[k-1]
// each w[i] strictly smaller than all its suffix
void duval(const string &s, vector<pii> &w) {
  for (int n = (int)s.size(), i = 0, j, k; i < n; ) {</pre>
    for (j = i + 1, k = i; j < n \& s[k] <= s[j]; j++)
      k = (s[k] < s[j] ? i : k + 1);
    // if (i < n / 2 \&\& j >= n / 2) {
    // for min cyclic shift, call duval(s + s)
    // then here s.substr(i, n / 2) is min cyclic shift
    for (; i <= k; i += j - k)</pre>
      w.pb(pii(i, j - k)); // s.substr(l, len)
 }
}
```

8.4 Suffix Array [cd67ea]

```
struct SuffixArray {
  vector<int> sa, lcp, rank; // lcp[i] is lcp of sa[i]
      and sa[i-1]
                             // sa[0] = s.size()
                             // character should be 1-
                                  hased
```

```
21
  SuffixArray(string& s, int lim=256) { // or
      basic_string<int>
    int n = s.size() + 1, k = 0, a, b;
    vector<int> x(n, 0), y(n), ws(max(n, lim));
    rank.assign(n, 0);
    for (int i = 0; i < n - 1; i++) x[i] = s[i];</pre>
    sa = lcp = y, iota(sa.begin(), sa.end(), 0);
for (int j = 0, p = 0; p < n; j = max(1, j * 2),</pre>
         lim = p) {
      p = j, iota(y.begin(), y.end(), n - j);
for (int i = 0; i < n; i++)</pre>
         if (sa[i] >= j) y[p++] = sa[i] - j;
      for (int &i : ws) i = 0;
      for (int i = 0; i < n; i++) ws[x[i]]++;</pre>
      for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];</pre>
      for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
       swap(x, y), p = 1, x[sa[0]] = 0;
      for(int i = 1; i < n; i++){</pre>
         a = sa[i - 1], b = sa[i];
        x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ?
              p - 1 : p++;
      }
    for (int i = 1; i < n; i++) rank[sa[i]] = i;</pre>
    for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
      for (k \&\& k--, j = sa[rank[i] - 1];
           s[i + k] == s[j + k]; k++);
};
8.5
     Suffix Automaton [016373]
// == a14210 ==
struct exSAM {
  const int CNUM = 26;
  // len: maxlength, link: fail link
  // LenSorted: topo order, cnt: occur
  vector<int> len, link, lenSorted, cnt;
  vector<vector<int>> next;
  int total = 0;
  int newnode() {
    return total++:
  void init(int n) { // total number of characters
    len.assign(2 * n, 0); link.assign(2 * n, 0);
    lenSorted.assign(2 * n, 0); cnt.assign(2 * n, 0);
    next.assign(2 * n, vector<int>(CNUM));
    newnode(), link[0] = -1;
// == c83c9c ==
  int insertSAM(int last, int c) { // 081739
    // not exSAM: cur = newnode(), p = Last
    int cur = next[last][c];
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
    if (p == -1) return link[cur] = 0, cur;
    int q = next[p][c];
    if (len[p] + 1 == len[q]) return link[cur] = q, cur
    int clone = newnode();
    for (int i = 0; i < CNUM; ++i)</pre>
      next[clone][i] = len[next[q][i]] ? next[q][i] :
           0;
    len[clone] = len[p] + 1;
    while (p != -1 && next[p][c] == q)
      next[p][c] = clone, p = link[p];
    link[link[cur] = clone] = link[q];
    link[q] = clone;
    return cur;
  void insert(const string &s) { // e47d43
    int cur = 0;
    for (auto ch : s) {
      int &nxt = next[cur][int(ch - 'a')];
      if (!nxt) nxt = newnode();
      cnt[cur = nxt] += 1;
// == 0a715a ==
```

void build() {

queue<int> q;

```
q.push(0);
    while (!q.empty()) {
      int cur = q.front();
      q.pop();
      for (int i = 0; i < CNUM; ++i)</pre>
        if (next[cur][i])
          q.push(insertSAM(cur, i));
    }
    vector<int> lc(total);
    for (int i = 1; i < total; ++i) ++lc[len[i]];</pre>
    partial_sum(iter(lc), lc.begin());
    for (int i = 1; i < total; ++i) lenSorted[--lc[len[</pre>
        i]]] = i;
  void solve() {
    for (int i = total - 2; i >= 0; --i)
      cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
  }
};
8.6 Z-value Algorithm [488d87]
// z[i] = max \ k \ s.t. \ s[0..k-1] = s[i..i+k-1]
// i.e. length of longest common prefix
// z[0] = 0
vector<int> z function(const string &s){
  int n = s.size();
  vector<int> z(n);
  for(int i = 1, l = 0, r = 0; i < n; i++){</pre>
    if(i <= r) z[i] = min(r - i + 1, z[i - 1]);</pre>
    while(i + z[i] < n && s[z[i]] == s[i + z[i]])
     z[i]++;
    if(i + z[i] - 1 > r)
      l = i, r = i + z[i] - 1;
  }
  return z;
}
8.7 Main Lorentz [fcfb8f]
struct Rep{ int minl, maxl, len; };
vector<Rep> rep; // 0-base
// p \in [minl, maxl] => s[p, p + i) = s[p + i, p + 2i)
void main_lorentz(const string &s, int sft = 0) {
  const int n = s.size();
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
  const string u = s.substr(0, nu), v = s.substr(nu),
        ru(u.rbegin(), u.rend()), rv(v.rbegin(), v.rend
            ());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const auto z1 = z_function(ru), z2 = z_function(v + '
      #' + u),
             z3 = z_{function}(ru + '#' + rv), z4 =
                  z_function(v);
  auto get_z = [](const vector<int> &z, int i) {
    return (0 <= i and i < (int)z.size()) ? z[i] : 0;</pre>
  auto add_rep = [&](bool left, int c, int l, int k1,
      int k2) {
    const int L = max(1, 1 - k2), R = min(1 - left, k1)
    if (L > R) return;
    if (left) rep.emplace_back(Rep({sft + c - R, sft +
        c - L, 1}));
    else rep.emplace_back(Rep({sft + c - R - l + 1, sft
         + c - L - 1 + 1, 1));
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int 1, k1, k2;
    if (cntr < nu) {</pre>
      1 = nu - cntr;
      k1 = get_z(z1, nu - cntr);
      k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= 1)
      add_rep(cntr < nu, cntr, 1, k1, k2);</pre>
  }
}
```

```
8.8 AC Automaton [f529e6]
```

```
const int SIGMA = 26;
struct AC_Automaton {
  // child: trie, next: automaton
  vector<vector<int>> child, next;
  vector<int> fail, cnt, ord;
  int total = 0;
  int newnode() {
    return total++;
  void init(int len) { // len >= 1 + total len
    child.assign(len, vector<int>(26, -1));
    next.assign(len, vector<int>(26, -1));
    fail.assign(len, -1); cnt.assign(len, 0);
    ord.clear();
    newnode();
  int input(string &s) {
    int cur = 0;
    for (char c : s) {
      if (child[cur][c - 'A'] == -1)
    child[cur][c - 'A'] = newnode();
      cur = child[cur][c - 'A'];
    return cur; // return the end node of string
  void make_fl() {
    queue<int> q;
    q.push(0), fail[0] = -1;
    while(!q.empty()) {
      int R = q.front();
      q.pop(); ord.pb(R);
      for (int i = 0; i < SIGMA; i++)</pre>
        if (child[R][i] != -1) {
          int X = next[R][i] = child[R][i], Z = fail[R
          while (Z != -1 && child[Z][i] == -1)
            Z = fail[Z];
          fail[X] = Z != -1 ? child[Z][i] : 0;
          q.push(X);
        else next[R][i] = R ? next[fail[R]][i] : 0;
    }
  void solve() {
    for (int i : ord | views::reverse)
      if (i) cnt[fail[i]] += cnt[i];
  }
};
      Palindrome Automaton [8a071b]
```

8.9

```
struct PalindromicTree {
 struct node {
    int nxt[26], fail, len; // num = depth of fail link
    int cnt, num; // cnt = occur, num = #pal_suffix of
        this node
    node(int l = 0) : nxt{},fail(0),len(1),cnt(0),num
        (0) {}
 vector<node> st; vector<int> s; int last, n;
 void init() {
    st.clear(); s.clear(); last = 1; n = 0;
    st.pb(0); st.pb(-1);
    st[0].fail = 1; s.pb(-1);
 int getFail(int x) {
    while (s[n - st[x].len - 1] != s[n]) x = st[x].fail
    return x;
 void add(int c) {
    s.pb(c -= 'a'); ++n;
    int cur = getFail(last);
    if (!st[cur].nxt[c]) {
      int now = SZ(st);
      st.pb(st[cur].len + 2);
      st[now].fail = st[getFail(st[cur].fail)].nxt[c];
     st[cur].nxt[c] = now;
     st[now].num = st[st[now].fail].num + 1;
    last = st[cur].nxt[c]; ++st[last].cnt;
```

```
void dpcnt() {
  for(int i = SZ(st) - 1; i >= 0; i--){
    auto nd = st[i];
    st[nd.fail].cnt += nd.cnt;
int size() { return (int)st.size() - 2; }
```

8.10 Palindrome Partition [c85c05]

```
// in PAM
/* node */ int dif = 0, slink = 0, g = 0;
vector<int> dp = {0};
if (!st[cur].nxt[c]) {
  st[now].dif = st[now].len - st[st[now].fail].len;
  if (st[now].dif == st[st[now].fail].dif)
    st[now].slink = st[st[now].fail].slink;
  else st[now].slink = st[now].fail;
dp.pb(0);
for (int x = last; x > 1; x = st[x].slink) {
  st[x].g = dp[n - st[st[x].slink].len - st[x].dif];
  if (st[x].dif == st[st[x].fail].dif)
    st[x].g = min(st[x].g, st[st[x].fail].g);
  dp[n] = min(dp[n], st[x].g + 1);
```

Misc 9

Cyclic Ternary Search [9017cc]

```
/* bool pred(int a, int b);
f(0) \sim f(n - 1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
while (r - 1 > 1) {
    int m = (1 + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else 1 = m;
  return pred(1, r % n) ? 1 : r % n;
```

9.2 Matroid

 $M=(E,\mathcal{I})$, where $\mathcal{I}\subseteq 2^E$ is nonempty, is a matroid if:

- If $S \in \mathcal{I}$ and $S' \subseteq S$, then $S' \in \mathcal{I}$.
- For $S_1, S_2 \in \mathcal{I}$ s.t. $|S_1| < |S_2|$, there exists $e \in S_2 \setminus S_1$ s.t. $S_1 \cup \{e\} \in \mathcal{I}$. Matroid intersection:

Start from $S=\emptyset.$ In each iteration, let

- $Y_1 = \{ x \not\in S \mid S \cup \{x\} \in \mathcal{I}_1 \}$
- $Y_2 = \{ x \notin S \mid S \cup \{x\} \in \mathcal{I}_2 \}$

If there exists $x \in Y_1 \cap Y_2$, insert x into S. Otherwise for each $x \in S, y \not \in S$, create edges

- $x \to y \text{ if } S \{x\} \cup \{y\} \in \mathcal{I}_1.$
- $y \to x$ if $S \{x\} \cup \{y\} \in \mathcal{I}_2$.

Find a *shortest* path (with BFS) starting from a vertex in Y_1 and ending at a vertex in Y_2 which doesn't pass through any other vertices in Y_2 , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if $x \in S$ and -w(x)if $x \not \in S$. Find the path with the minimum number of edges among all minimum length paths and alternate it

9.3 Simulate Annealing [ff826c]

```
ld anneal() {
  mt19937 rnd_engine(seed);
  uniform real distribution<ld> rnd(0, 1);
  const ld dT = 0.001;
  // Argument p
  ld S_cur = calc(p), S_best = S_cur;
  for (1d T = 2000; T > eps; T -= dT) {
    // Modify p to p_prime
    const ld S_prime = calc(p_prime);
    const ld delta_c = S_prime - S_cur;
    ld prob = min((ld)1, exp(-delta_c / T));
    if (rnd(rnd_engine) <= prob)</pre>
      S_cur = S_prime, p = p_prime;
    if (S_prime < S_best) // find min</pre>
```

```
S_best = S_prime, p_best = p_prime;
return S_best;
```

9.4 Binary Search On Fraction [f6b9ec]

```
struct Q {
  11 p, q;
  Q go(Q b, 11 d) { return {p + b.p * d, q + b.q * d};
// returns smallest p/q in [lo, hi] such that
// pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N, auto &&pred) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step /= 2 : step *=
       2);)
     if (Q mid = hi.go(lo, len + step);
         else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
```

9.5 Min Plus Convolution [09b5c3]

```
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution(vector<int> &a, vector
    <int> &b) {
  int n = SZ(a), m = SZ(b);
  vector<int> c(n + m - 1, INF);
  auto dc = [&](auto Y, int 1, int r, int jl, int jr) {
    if (1 > r) return;
    int mid = (1 + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)
  if (int i = mid - j; i >= 0 && i < n)</pre>
        if (best > a[i] + b[j])
          best = a[i] + b[j], from = j;
    Y(Y, 1, mid - 1, jl, from), Y(Y, mid + 1, r, from,
        jr);
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
```

9.6 SMAWK [a2a4ce]

```
// For all 2x2 submatrix:
// If M[1][0] < M[1][1], M[0][0] < M[0][1]
// If M[1][0] == M[1][1], M[0][0] <= M[0][1]
// M[i][ans_i] is the best value in the i-th row
// select(int r, int u, int v) return true if f(r, v)
    is better than f(r, u)
vector<int> smawk(int N, int M, auto &&select) {
  auto dc = [&](auto self, const vector<int> &r, const
      vector<int> &c) {
    if (r.empty()) return vector<int>{};
    const int n = SZ(r); vector<int> ans(n), nr, nc;
    for (int i : c) {
      while (!nc.empty() &&
          select(r[nc.size() - 1], nc.back(), i))
        nc.pop_back();
      if (int(nc.size()) < n) nc.push_back(i);</pre>
    for (int i = 1; i < n; i += 2) nr.push_back(r[i]);</pre>
    const auto na = self(self, nr, nc);
    for (int i = 1; i < n; i += 2) ans[i] = na[i >> 1];
    for (int i = 0, j = 0; i < n; i += 2) {
      ans[i] = nc[j];
      const int end = i + 1 == n ? nc.back() : ans[i +
          1];
      while (nc[j] != end)
        if (select(r[i], ans[i], nc[++j])) ans[i] = nc[
```

```
return ans;
vector<int> R(N), C(M); iota(iter(R), 0), iota(iter(C)
return dc(dc, R, C);
```

9.7 Golden Ratio Search [ce06a8]

```
ld goldenRatioSearch(ld a, ld b, auto &&f) {
  ld r = (sqrt(5)-1)/2, eps = 1e-7;
ld x1 = b - r*(b-a), x2 = a + r*(b-a);
ld f1 = f(x1), f2 = f(x2);
  while (b-a > eps)
     if (f1 < f2) { //change to > to find maximum
        b = x2; x2 = x1; f2 = f1;
        x1 = b - r*(b-a); f1 = f(x1);
     } else {
        a = x1; x1 = x2; f1 = f2; x2 = a + r*(b-a); f2 = f(x2);
  return a;
}
```

10 **Notes** 10.1 Geometry **Rotation Matrix**

$$\begin{pmatrix}
\cos\theta & -\sin\theta \\
\sin\theta & \cos\theta
\end{pmatrix}$$

- rotate 90° : $(x,y) \rightarrow (-y,x)$
- rotate -90° : $(x,y) \rightarrow (y,-x)$

Triangles

}

Side lengths: a, b, c

Semiperimeter:
$$p = \frac{a+b+c}{2}$$

Area:
$$A = \sqrt{p(p-a)(p-b)(p-c)}$$

Circumradius: $R = \frac{abc}{4}$

Inradius: $r = \frac{A}{}$

Length of median (divides triangle into two equal-area triangles): $m_a =$ $\frac{1}{2}\sqrt{2b^2+2c^2-a^2}$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc\left(1-\left(\frac{a}{b+c}\right)^2\right)}$$

Law of sines: $\frac{\sin\alpha}{a} = \frac{\sin\beta}{b} = \frac{\sin\gamma}{c} = \frac{1}{2R}$ Law of cosines: $a^2 = b^2 + c^2 - 2bc\cos\alpha$ Law of tangents: $\frac{a+b}{a-b} = \frac{\tan\frac{\alpha+\beta}{2}}{\tan\frac{\alpha-\beta}{2}}$

Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac + bd, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$. Spherical coordinates

$$\begin{array}{ll} x = r \sin \theta \cos \phi & r = \sqrt{x^2 + y^2 + z^2} \\ y = r \sin \theta \sin \phi & \theta = \mathrm{acos}(z/\sqrt{x^2 + y^2 + z^2}) \\ z = r \cos \theta & \phi = \mathrm{atan2}(y,x) \end{array}$$

Green's Theorem

$$\begin{split} \iint_D \left(\frac{\partial Q}{\partial x} - \frac{\partial P}{\partial y} \right) dx dy &= \oint_{L^+} (P dx + Q dy) \\ \text{Area} &= \frac{1}{2} \oint_{\mathbb{R}} x \ dy - y \ dx \end{split}$$

Circular sector:

$$\begin{split} x &= x_0 + r \cos \theta \\ y &= y_0 + r \sin \theta \\ A &= r \int_{\alpha}^{\beta} (x_0 + \cos \theta) \cos \theta + (y_0 + \sin \theta) \sin \theta \, d\theta \\ &= r (r\theta + x_0 \sin \theta - y_0 \cos \theta)|_{\alpha}^{\beta} \end{split}$$

· Centroid:

$$\bar{x} = \frac{1}{2A} \int_C x^2 dy \quad \bar{y} = -\frac{1}{2A} \int_C y^2 dx$$

Point-Line Duality

$$p = (a, b) \leftrightarrow p^* : y = ax - b$$

- $p \in l \iff l^* \in p^*$
- p_1, p_2, p_3 are collinear $\iff p_1^*, p_2^*, p_3^*$ intersect at a point
- p lies above $l \iff l^*$ lies above p
- lower convex hull \leftrightarrow upper envelope

10.2 Trigonometry

$$\sinh x = \frac{1}{2}(e^x - e^{-x}) \qquad \qquad \cosh x = \frac{1}{2}(e^x + e^{-x})$$

$$\sin n\pi = 0 \qquad \qquad \cos n\pi = (-1)^n$$

$$\sin(\alpha + \beta) = \sin\alpha\cos\beta + \cos\alpha\sin\beta$$

$$\cos(\alpha + \beta) = \cos\alpha\cos\beta - \sin\alpha\sin\beta$$

$$\sin(2\alpha) = 2\cos\alpha\sin\alpha$$

$$\cos(2\alpha) = \cos^2 \alpha - \sin^2 \alpha = 2\cos^2 \alpha - 1 = 1 - 2\sin^2 \alpha$$

$$\tan(\alpha + \beta) = \frac{\tan \alpha + \tan \beta}{1 - \tan \alpha \tan \beta}$$

$$\sin \alpha + \sin \beta = 2 \sin \frac{\alpha + \beta}{2} \cos \frac{\alpha - \beta}{2}$$

$$\cos \alpha + \cos \beta = 2 \cos \frac{\alpha + \beta}{2} \cos \frac{\alpha - \beta}{2}$$

$$\sin \alpha \sin \beta = \frac{1}{2}(\cos(\alpha - \beta) - \cos(\alpha + \beta))$$

$$\sin \alpha \cos \beta = \frac{1}{2}(\sin(\alpha + \beta) + \sin(\alpha - \beta))$$

$$\cos \alpha \sin \beta = \frac{1}{2}(\sin(\alpha + \beta) - \sin(\alpha - \beta))$$
$$\cos \alpha \cos \beta = \frac{1}{2}(\cos(\alpha - \beta) + \cos(\alpha + \beta))$$

$$(V+W)\tan(\alpha-\beta)/2 = (V-W)\tan(\alpha+\beta)/2$$

where V, W are lengths of sides opposite angles α, β .

$$a\cos x + b\sin x = r\cos(x - \phi)$$

$$a\sin x + b\cos x = r\sin(x+\phi)$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

10.3 Calculus

Integration by parts:

$$\int_{a}^{b} f(x)g(x)dx = [F(x)g(x)]_{a}^{b} - \int_{a}^{b} F(x)g'(x)dx$$

$$\frac{d}{dx}\arcsin x = \frac{1}{\sqrt{1-x^2}} \qquad \frac{d}{dx}\arccos x = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}\tan x = 1 + \tan^2 x \qquad \frac{d}{dx}\arctan x = \frac{1}{1 + x^2}$$

$$\frac{1}{dx} \arctan x = \frac{1}{1 + x^2}$$

$$\int \tan ax = -\frac{\ln|\cos ax|}{a} \qquad \int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2}$$

$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2} \operatorname{erf}(x) \qquad \int x e^{ax} = \frac{e^{ax}}{a^2} (ax - 1)$$

$$\int \sin^2(x) = \frac{x}{2} - \frac{1}{4}\sin 2x \qquad \int \sin^3 x = \frac{1}{12}\cos 3x - \frac{3}{4}\cos x$$

$$\int \cos^2(x) = \frac{x}{2} + \frac{1}{4} \sin 2x \qquad \int \cos^3 x = \frac{1}{12} \sin 3x + \frac{3}{4} \sin x$$

$$\int x \sin x = \sin x - x \cos x \qquad \int x \cos x = \cos x + x \sin x$$

$$\int xe^x = e^x(x-1) \qquad \int x^2e^x = e^x(x^2 - 2x + 2)$$

$$\int x^2 \sin x = 2x \sin x - (x^2 - 2) \cos x$$

$$\int x^2 \cos x = 2x \cos x + (x^2 - 2) \sin x$$

$$\int e^x \sin x = \frac{1}{2} e^x (\sin x - \cos x)$$

$$\int e^x \cos x = \frac{1}{2}e^x(\sin x + \cos x)$$

$$\int xe^x \sin x = \frac{1}{2}e^x(x\sin x - x\cos x + \cos x)$$

$$\int xe^x \cos x = \frac{1}{2}e^x(x\sin x + x\cos x - \sin x)$$

10.4 Sum & Series

$$\begin{split} c^a + c^{a+1} + \dots + c^b &= \frac{c^{b+1} - c^a}{c-1}, c \neq 1 \\ 1 + 2 + 3 + \dots + n &= \frac{n(n+1)}{2} \\ 1^2 + 2^2 + 3^2 + \dots + n^2 &= \frac{n(2n+1)(n+1)}{6} \\ 1^3 + 2^3 + 3^3 + \dots + n^3 &= \frac{n^2(n+1)^2}{4} \\ 1^4 + 2^4 + 3^4 + \dots + n^4 &= \frac{n(n+1)(2n+1)(3n^2 + 3n - 1)}{30} \\ e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty) \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \le 1) \\ \sqrt{1+x} &= 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1) \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty) \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty) \end{split}$$

10.5 Misc

· Cramer's rule

$$ax + by = e$$

$$cx + dy = f \Rightarrow x = \frac{ed - bf}{ad - bc}$$

$$y = \frac{af - ec}{ad - bc}$$

Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

· Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.
- BEST theorem: the number of eulerian circuits in a directed graph is $|\mathsf{det}(\tilde{L}_{ww})| \cdot \prod_{v \in V} (\mathsf{deg}(v) - 1)!$.
- Tutte's Matrix

Let D be a n imes n matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on \widetilde{G} .

- Cayley's Formula
 - Given a degree sequence d_1, d_2, \dots, d_n for each $\emph{labeled}$ vertices, there are $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
 - Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.
- Erdős–Gallai theorem

A sequence of nonnegative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on \boldsymbol{n} vertices if and only

if
$$d_1+\cdots+d_n$$
 is even and $\sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k)$ holds

for every $1 \le k \le n$.

Gale-Ryser theorem

A pair of sequences of nonnegative integers $a_1 \geq \cdots \geq a_n$ and b_1, \ldots, b_n

is bigraphic if and only if
$$\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$$
 and $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$

holds for every 1 < k < n.

Fulkerson-Chen-Anstee theorem

A sequence $(a_1, b_1), \ldots, (a_n, b_n)$ of nonnegative integer pairs with $a_1 \geq$

$$\cdots \geq a_n$$
 is digraphic if and only if $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$ and $\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i, k_i)$

$$1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

For simple polygon, when points are all integer, we have $A=\#\{\text{lattice points in the interior}\}+\frac{\#\{\text{lattice points on the boundary}\}}{2}-1.$

- · Möbius inversion formula
- $\mu(d)=(-1)^k$ if n is the product of k distinct primes, 0 if $p^2\mid n$ $f(n)=\sum_{d\mid n}g(d)\Leftrightarrow g(n)=\sum_{d\mid n}\mu(d)f(\frac{n}{d})$

- $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- · Spherical cap
- A portion of a sphere cut off by a plane.
- r: sphere radius, a: radius of the base of the cap, h: height of the cap,
- Volume = $\pi h^2 (3r h)/3 = \pi h(3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 h^2)$
- Area $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta).$
- · Lagrange multiplier
 - Optimize $f(x_1,\ldots,x_n)$ when k constraints $g_i(x_1,\ldots,x_n)=0$.
 - Lagrangian function
 - $\mathcal{L}(x_1,\ldots,x_n,\lambda_1,\ldots,\lambda_k) = f(x_1,\ldots,x_n) \sum_{i=1}^k \lambda_i g_i(x_1,\ldots,x_n).$
 - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- · Nearest points of two skew lines
 - Line 1 : ${m v}_1 = {m p}_1 + t_1 {m d}_1$
 - Line 2 : ${m v}_2 = {m p}_2 + t_2 {m d}_2$
 - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$

 - $\begin{array}{l} \textbf{-} \ \, \boldsymbol{n} = \boldsymbol{a}_1 \wedge \boldsymbol{\omega}_2 \\ \textbf{-} \ \, \boldsymbol{n}_1 = \boldsymbol{d}_1 \times \boldsymbol{n} \\ \textbf{-} \ \, \boldsymbol{n}_2 = \boldsymbol{d}_2 \times \boldsymbol{n} \\ \textbf{-} \ \, \boldsymbol{c}_1 = \boldsymbol{p}_1 + \frac{(\boldsymbol{p}_2 \boldsymbol{p}_1) \cdot \boldsymbol{n}_2}{d_1 \cdot \boldsymbol{n}_2} \boldsymbol{d}_1 \\ \textbf{-} \ \, \boldsymbol{c}_2 = \boldsymbol{p}_2 + \frac{(\boldsymbol{p}_1 \boldsymbol{p}_2) \cdot \boldsymbol{n}_1}{d_2 \cdot \boldsymbol{n}_1} \boldsymbol{d}_2 \end{array}$

Bernoulli numbers
$$B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0$$

$$\sum_{j=0}^m {m+1 \choose j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}.$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k^+ n^{m+1-k}$$

• Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^n$$

$$x^n = \sum_{i=0}^{n} S(n,i)(x)_i$$
 • Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left(x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$
 • Catalan numbers
$$C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$$

$$C_n^{(k)} = \frac{1}{(k-1)n+1} {kn \choose n}$$
$$C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$\begin{aligned} &\pi(j) \geq j, k \ j \text{:s s.t. } \pi(j) > j. \\ &E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k) \\ &E(n,0) = E(n,n-1) = 1 \\ &E(n,k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n \end{aligned}$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{i} {n+1 \choose i} (k+1-j)^{n}$$

10.6 Number

• Some prime numbers:

12721, 13331, 14341, 75577, 123457, 222557, 556679, 999983, 1097774749, 1076767633, 100102021, 999997771, 1001010013, 1000512343, 987654361, 999991231, 999888733, 98789101, 987777733, 999991921, 1010101333, 1010102101, 1000000000039 100000000000037, 2305843009213693951, 4611686018427387847, 9223372036854775783, 18446744073709551557

• Number of paritions of *n*: *n* | 2 3 4 5 6 7 8 9 20 30 40 50 100 p(n) 2 3 5 7 11 15 22 30 627 5604 4e4 2e5 2e8

Maximum number of divisors:

 $n \mid$ 100 1e3 1e6 1e9 1e12 1e15 $\overline{d(i)}$ 12 32 240 1344 6720 26880 103680

- n 12345678 9 10 11 12 13 14 15 $\frac{\binom{2n}{n}}{2}$ 6 20 70 252 924 3432 12870 48620 184756 7e5 2e6 1e7 4e7 1.5e8

• Fibonacci numbers: $\frac{n}{F_n} \begin{vmatrix} 1 & 2 & 3 & 4 & 5 & 31 & 45 & 88 \\ 1 & 1 & 1 & 3 & 5 & 8 & 1346269 & 1e9 & 1e18 \end{vmatrix}$

