Week13- Write-up

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Challenge: Super Secure Letter

Objective

The challenge presents an encrypted message (ciphertext) that we need to decrypt. The encrypted message is generated from a flag using an XOR-based encryption method combined with PRNG.

Obective is to recover the plaintext (flag) from the given ciphertext by reverse engineering the encryption logic used in the binary.

Approach

I analyzed the provided binary in Binary Ninja and identified the encryption logic. Here's how the encryption works:

1. The Seed for PRNG

The binary seeds the pseudo-random number generator (PRNG) using the following formula:

srand(time(nullptr) * time(nullptr));

- time(nullptr) fetches the current Unix timestamp.
- The seed for srand() is the square of the timestamp.

The XOR Encryption

After seeding, the binary loops over the flag bytes and XORs each byte with the least significant byte (LSB) of the PRNG output generated using rand():

```
for (int i = 0; i < strlen(flag); i++) {
printf("%02x", rand() & 0xFF ^ flag[i]);
}</pre>
```

- rand(): Produces the next random number.
- rand() & 0xFF: Extracts the least significant byte.
- flag[i]: The current byte of the flag.
- XOR (^) is applied between the PRNG output and the flag byte.

Decryption Plan

The XOR operation is symmetric, so the encryption process can be reversed:

Original Byte = Ciphertext Byte ^ PRNG Output

To recover the flag:

- Brute-force potential seeds by squaring timestamps near the current time.
- Replicate the PRNG sequence using libc.srand(seed) and libc.rand().
- XOR the PRNG output with the ciphertext bytes.

Here is the Python code I used to solve the challenge:

```
~/Super_Secret_Letter.py - Mousepad
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 2 import ctypes, re, time
 4 # Load libc for C's rand() and srand()
 5 libc = ctypes.CDLL("libc.so.6")
 6 libc.srand.argtypes, libc.rand.restype = [ctypes.c_uint], ctypes.c_int
 9 def decrypt(ciphertext, seed):
    libc.srand(seed)
return bytes(c ^ (libc.rand() & 0×FF) for c in ciphertext)
11
12
13 # Connect to server and extract ciphertext
14 p = remote('offsec-chalbroker.osiris.cyber.nyu.edu', 1517)
15 p.sendlineafter(b'abc123): ', b'vc2499')
16 response = p.recvall(timeout=2).decode()
17 ciphertext = bytes.fromhex(re.search(r'[a-fA-F0-9]{64,}', response).group(0))
18 p.close()
19
20 # Brute-force squared seeds
21 current_time = int(time.time())
22 for t in range(current_time - 3600, current_time + 3600):
24
           plaintext = decrypt(ciphertext, t * t).decode('utf-8')
           if "flag{" in plaintext:
                    (f"[+] Got the Flag!\nSeed: {t * t}\nPlaintext: {plaintext}")
26
               exit()
27
28
29
30
31 print("[-] Flag not found.")
```

I ran the script and the code successfully recovered the flag:

flag{p3rh4p5_n07_50000_53cur3:(_8e3da4fc176ccfdb}

Challenge: Pseudo Rand

Objective

The challenge Pseudo Rand requires me to guess a "random" number generated by the server. The number is generated using a weak PRNG (rand()) seeded with the current time plus a fixed offset.

My objective was to predict the "random" number generated by the server based on the weak seeding mechanism and retrieve the flag

Analyzing the Binary

I started by opening the provided binary in Binary Ninja. I navigated to the main() function, where I noticed the following logic:

```
srand(time(nullptr) + 0x19);
```

int generated number = rand();

This immediately stood out. Here's what I understood:

- The random number generator (rand()) is seeded using the current time (time(nullptr)) plus an offset (+0x19).
- Since rand() is deterministic, I could replicate this locally if I knew the seed value.

```
      000012fa
      e851feffff
      call time

      000012ff
      83c019
      add eax, 0x19

      00001302
      89c7
      mov edi, eax

      00001304
      e817feffff
      call srand

      00001309
      e882feffff
      call rand
```

To predict the server's random number, I decided to use Python. I loaded the libc.so.6 library using ctypes because it contains the srand() and rand() functions, mirroring the C behavior.

Here's the step-by-step logic I implemented:

- 1. Compute the seed as current time + 0x19.
- 2. Call srand() with this seed to initialize the PRNG.
- 3. Use rand() to generate the first random number, which would match the server's output.

Script I wrote:

```
~/Psudo_Rand.py - Mousepad
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 1 from pwn import *
2 import ctypes
 4 # Connect to the remote server
 5 server = remote('offsec-chalbroker.osiris.cyber.nyu.edu', 1514)
 7 # Step 1: Send the NetID
8 netid = b'vc2499'
 9 server.sendlineafter(b'abc123): ', netid)
11 # Step 2: Predict the random number
12 libc - ctypes.CDLL("libc.so.6")  # Load the C standard library
13 current_seed = int(time.time()) + 0×19  # Compute the seed based on the challenge logic
14 libc.srand(current_seed)
                                    # Seed the PRNG
15 predicted_number = libc.rand()  # Generate the predicted random number
17 # Step 3: Send the predicted number to the server
18 server.sendlineafter(b"Please wait a moment...", str(predicted_number).encode())
19
20 # Step 4: Switch to interactive mode to receive the flag
21 server.interactive()
22
```

I ran the script and it retrieved the flag:

flag{l00ks_l1k3_th4t_seed_w4snt_gr34t!_8c728b5b0604c5d0}

```
(kali@ kali)-[~]
    $ python3 Psudo_Rand.py
[+] Opening connection to offsec-chalbroker.osiris.cyber.nyu.edu on port 1514: Done
[*] Switching to interactive mode

Can you guess my number?

I guess it wasn't that random!
    Here's your flag, friend: flag{l00ks_l1k3_th4t_seed_w4snt_gr34t!_8c728b5b0604c5d0}

[*] Got EOF while reading in interactive
```

Challenge: RSA_1

Objective

The goal of this challenge was to decrypt a ciphertext that was encrypted using RSA with a small public exponent e=5. By exploiting the weakness caused by the small exponent, I was able to recover the plaintext directly.

Solution Steps

Connecting to the Server

I connected to the server using pwntools and sent my NetID to retrieve the challenge parameters:

- e: Public exponent (fixed at 5)
- n: Modulus
- c: Ciphertext

```
Please wait a moment while we generate your ciphertext ...

e = 5
n = 9236124135376191527559713985703603262247125684837692826894734457523373879711990264424395588393214001622789984239540692795946561137470540356751
1721662486095223922000565903570881324256217704320709904363008579844817852956588557453497298977100093235943385123701242193862885751157660301113250
218452613094126099853808500724731097175000242495469002434650652902190190909141123249814422730257317639621691885771308881852687764941715916429779261
94083398577002881058727020999415880006622797393994442929449849683554935468158740390429323172034070725277972039474518553166858083757159422489983162059
5519302216038429110062837258670291216653039591474275573490524718193146424628120610595956610804421571183223766196487518745945283932249731659742504
6485058768200381825119884347125265797831398328744087393193032112809241558137567973162734584294895777272770889720936723870891021180037990885505277930
745584036110444611577761924893982239453106969536688468765563892066666325655585899036666325655585859253735764047581456062398557873518778
430708911791503025862745900255790833267144951643334628448610160038356407877392193095600153217379376333970787070380334483094715925725741366149661070
23916734418180321882484933674584926013009520025906442147963135639421
c = 73246603240320346337487822014005910509105151209774800158862190801231052519913105485896352957396810155545657830442308991488611233518011542876362
5622450178843121137339397928536743147462247232800498601734297522121125039472891391445471406605630693905465070823315395217821827114090818837556493058255706390509165151209774840015886219080123106251991310548589655699305466070823315395217821827114090818837556493058259
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7538282516092746455818141947233280049086017342975522121125090473459150901575834957526308107172345312059902575338915916400
3580156556784561222922978954161857042797180189729342672547322842354
```

RSA encryption is performed as:

c=m^e mod n

When e is small (in this case e=5) and m^e < n, the ciphertext c becomes simply m^e without wrapping due to the modulus n.

Thus, to recover the plaintext m, I computed the **5th root** of c.

I wrote a script using pwntools to:

- 1. Connect to the server and retrieve e,n,c.
- 2. Compute the 5th root of ccc using gmpy2.iroot.
- 3. Convert the result into a readable flag.

Script Used:

I ran the script and it retrieved the flag:

flag{n0_f4ct0r1ng_r3qu1r3d!_0c53cc2c0ebd78a6}

Challenge: RSA_2

Objective

The goal of this challenge was to decrypt a message that was encrypted using RSA with the **same modulus n** but two different public exponents e1 and e2. Using the vulnerability in this setup, I recovered the original plaintext message.

Solution Steps

Analyzing the Problem

I connected to the server using nc (netcat) and observed the following:

- Two public exponents e1 and e2.
- A shared modulus n.
- Two ciphertexts c1 and c2, which were encrypted using the same modulus n but with different exponents.

```
Please wait a moment while we generate your ciphertexts ...

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```

The challenge description suggested that this setup is vulnerable to the **Common Modulus Attack** in RSA.

In RSA, if two ciphertexts c1 and c2 are encrypted with the same modulus n but different public exponents e1 and e2, we can use **Bezout's Identity** to recover the plaintext m.

Bezout's Identity states that for integers e1 and e2, there exist coefficients a and b such that:

 $a \cdot e1 + b \cdot e2 = 1$

This property allows us to combine c1and c2 to recover m as follows:

m=(c1^a·c2^b) mod n

If b is negative, we use the **modular inverse** of c2

To solve the challenge, I wrote the following Python script using pwntools and gmpy2:

- which connected to the server and retrieved the values for e1,n,c1,e2,n,c2.
- Using gmpy2.gcdext, I calculated the coefficients a and b.
- If b was negative, I computed the modular inverse of c2 using gmpy2.invert
- I combined the ciphertexts c1 and c2 using the formula: m=(c1^a·c2^b) mod n
- Finally, I converted the decrypted integer m into a hexadecimal string and decoded it into ASCII to extract the flag.

Script used:

```
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```

I ran the script and received the flag:

flag{n1c3_j0b_br34k1ng_T3xtB00k_RSA!_c82781d6da68ab8d}