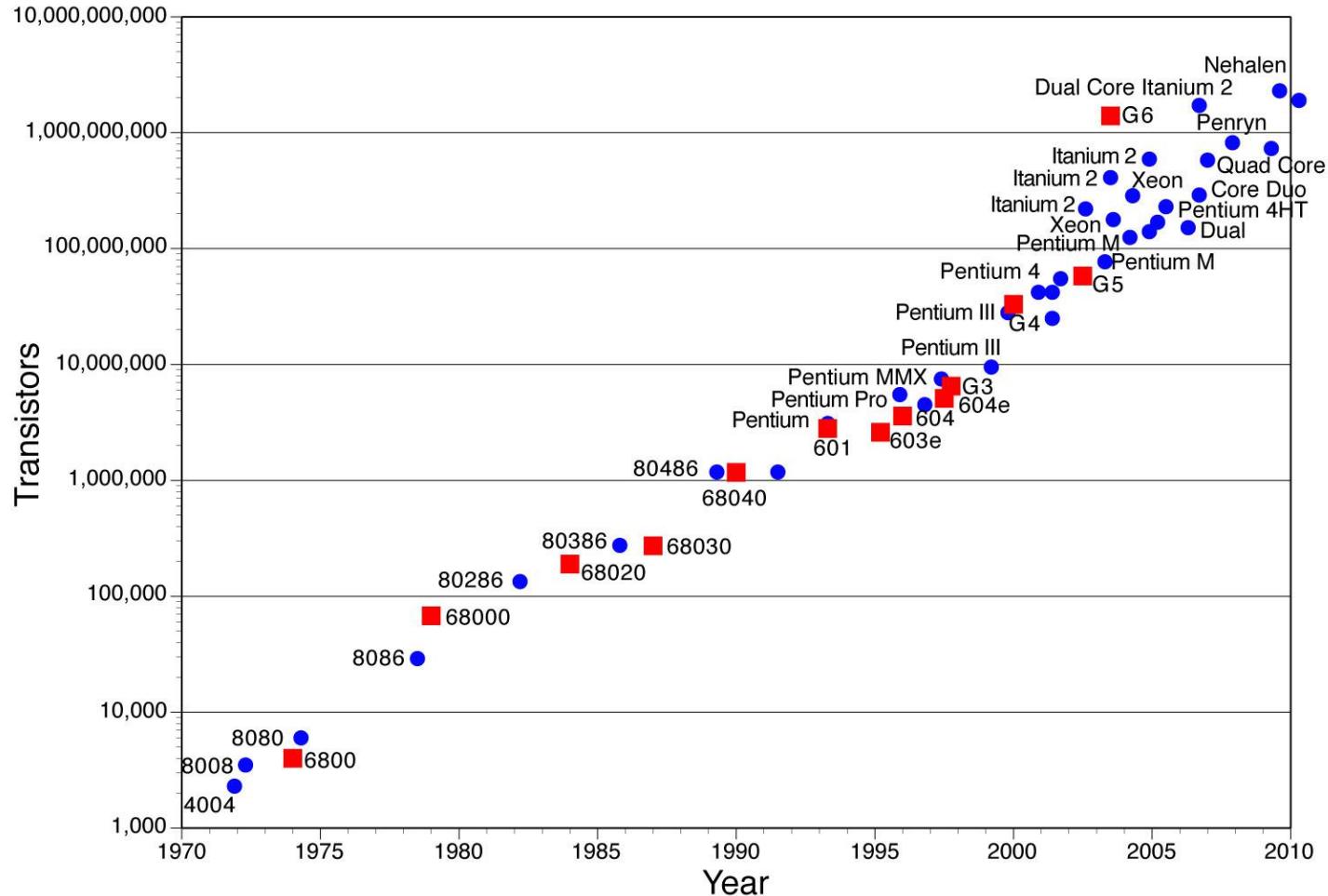
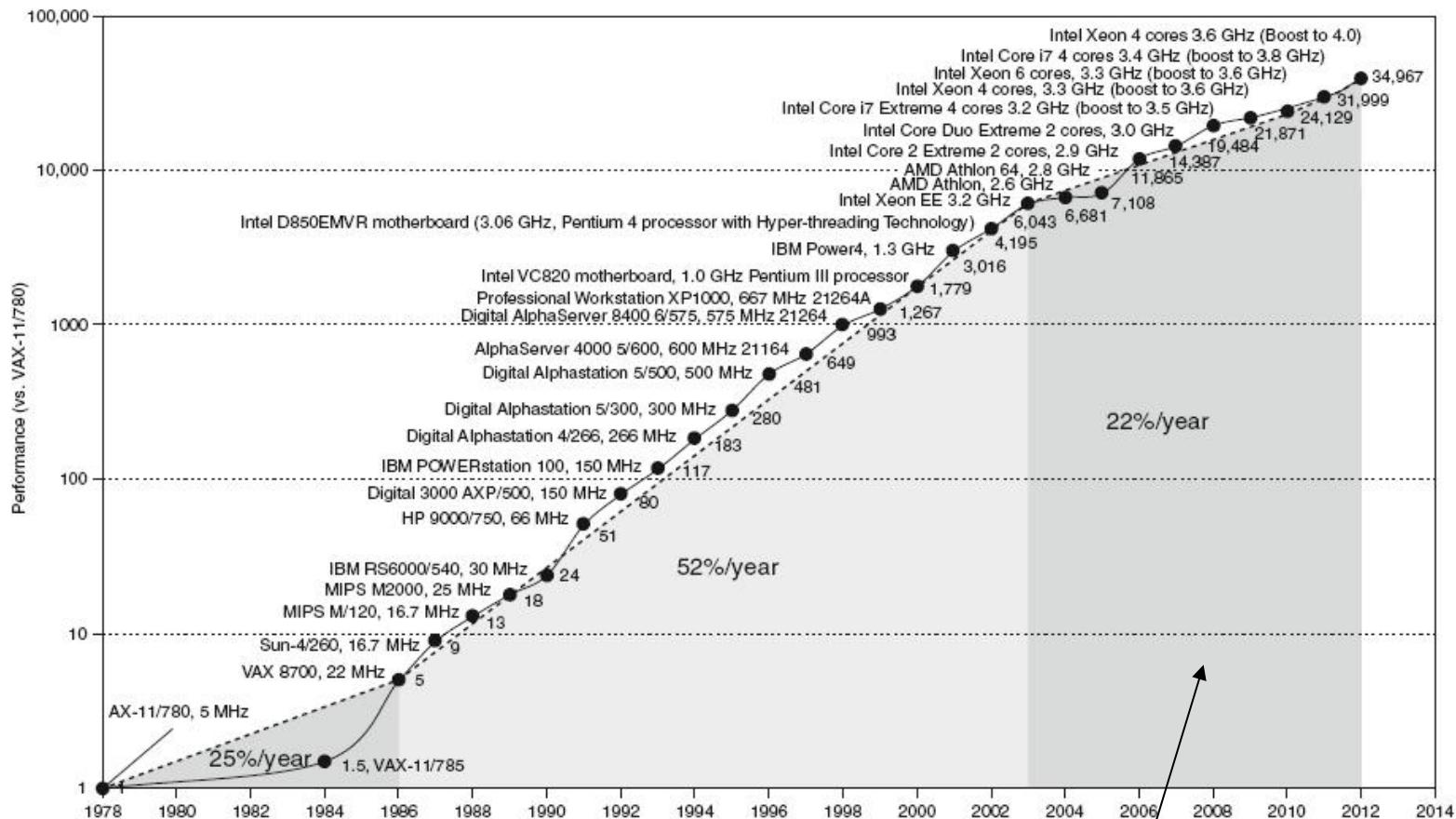


Overview of Prerequisites

Moore's Law

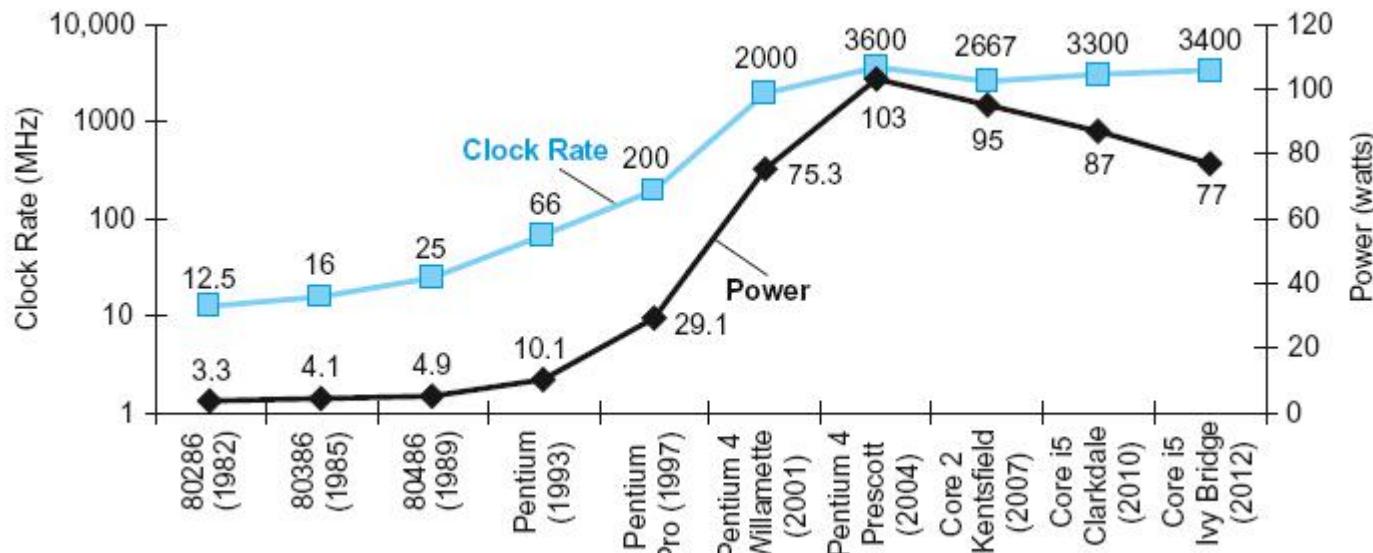


Uniprocessor Performance



Constrained by power, instruction-level parallelism,
memory latency

Power Trends



- In CMOS IC technology

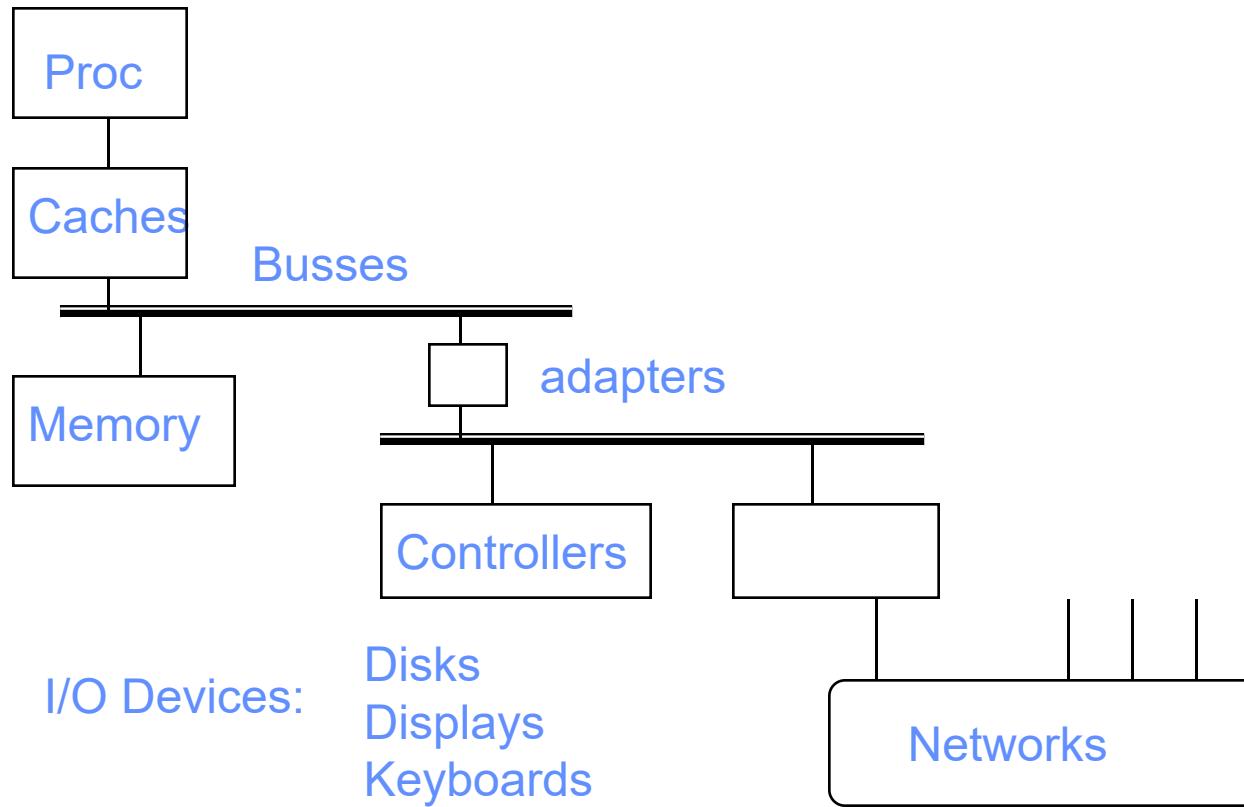
$$\text{Power} = \text{Capacitive load} \times \text{Voltage}^2 \times \text{Frequency}$$

×30

5V → 1V

×1000

Computer System Components



- All have interfaces & organizations
 - Bus & Bus Protocol is key to composition
- => peripheral hierarchy

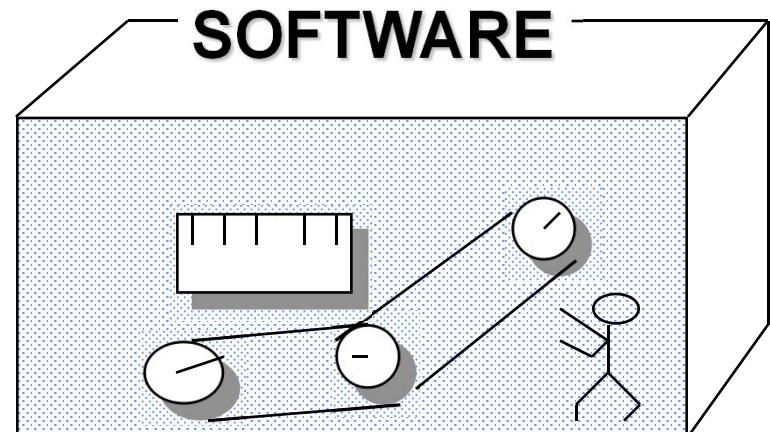
Instruction Set Architecture

... the attributes of a [computing] system as seen by the programmer, i.e. the conceptual structure and functional behavior, as distinct from the organization of the data flows and controls the logic design, and the physical implementation.

- Amdahl, Blaaw, and

Brooks, 1964

- Organization of Programmable Storage
- Data Types & Data Structures:
 Encodings & Representations
- Instruction Formats
- Instruction (or Operation Code) Set
- Modes of Addressing and Accessing Data Items and Instructions
- Exceptional Conditions



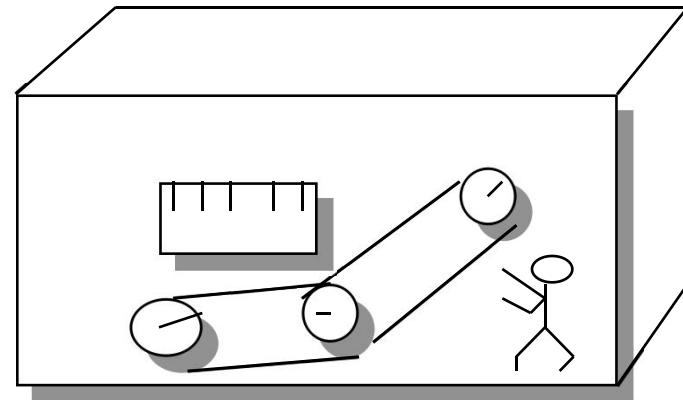
Organization

- Capabilities & Performance Characteristics of Principal Functional Units
 - (e.g., Registers, ALU, Shifters, Logic Units, ...)
- Ways in which these components are interconnected
- Information flows between components
- Logic and means by which such information flow is controlled.
- Choreography of FUs to realize the ISA
- Register Transfer Level (RTL) Description

Logic Designer's View

ISA Level

FUs & Interconnect



Review: Basic ISA Classes

Accumulator:

1 address	add A	$acc \leftarrow acc + mem[A]$
1+x address	addx A	$acc \leftarrow acc + mem[A + x]$

Stack:

0 address	add	$tos \leftarrow tos + next$
-----------	-----	-----------------------------

General Purpose Register:

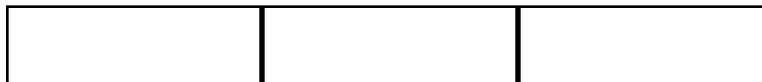
2 address	add A B	$EA(A) \leftarrow EA(A) + EA(B)$
3 address	add A B C	$EA(A) \leftarrow EA(B) + EA(C)$

Load/Store:

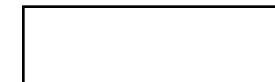
3 address	add Ra Rb Rc	$Ra \leftarrow Rb + Rc$
	load Ra Rb	$Ra \leftarrow mem[Rb]$
	store Ra Rb	$mem[Rb] \leftarrow Ra$

Instruction Formats

Variable:

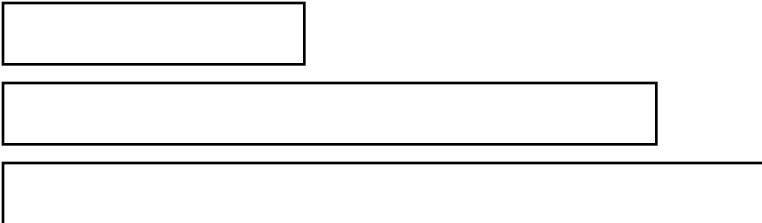


...



Fixed:

Hybrid:

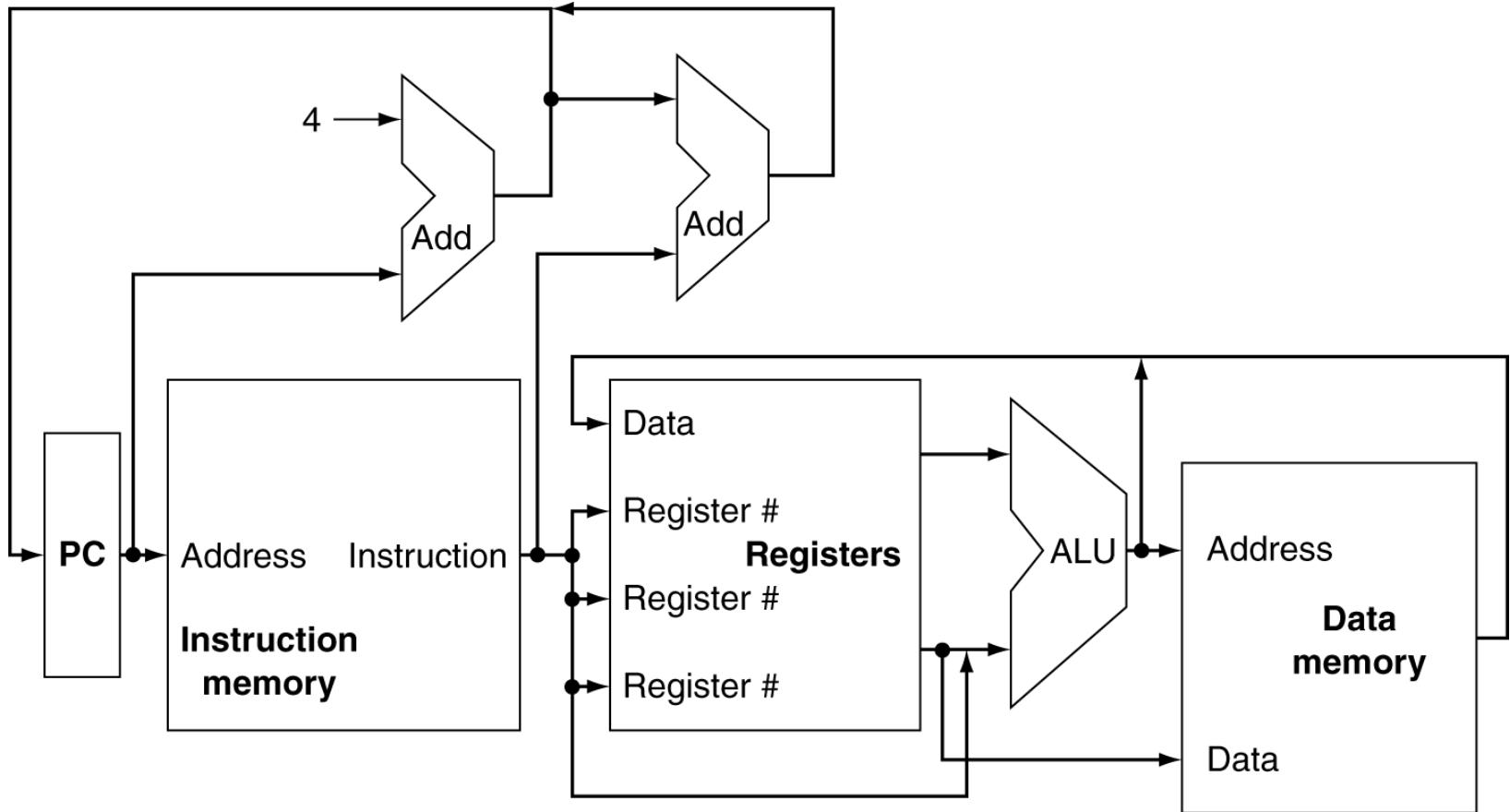


- Addressing modes
 - each operand requires address specifier => variable format
- code size => variable length instructions
- performance => fixed length instructions
 - simple decoding, predictable operations
- With load/store instruction arch, only one memory address and few addressing modes
- => simple format, address mode given by opcode

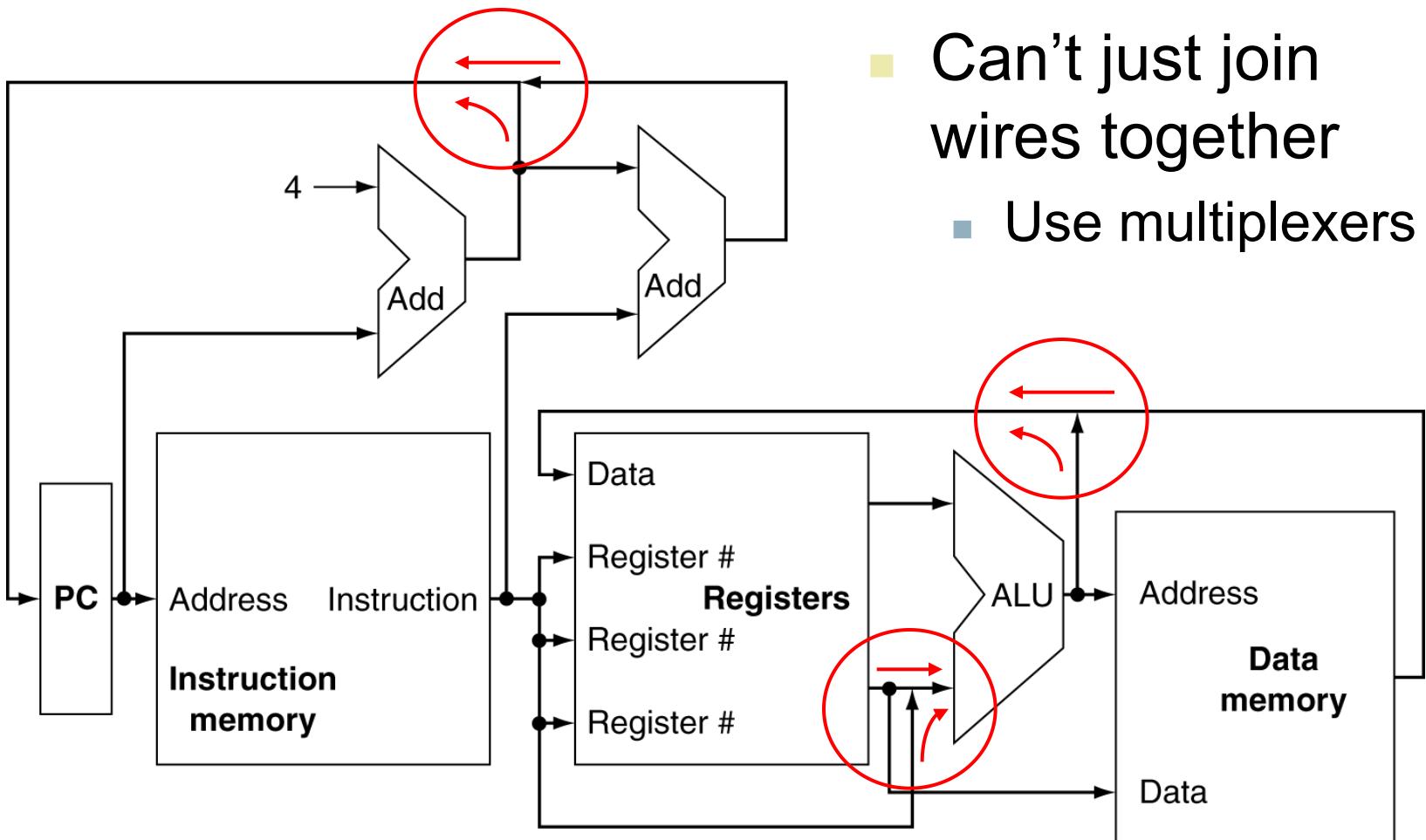
Instruction Execution

- PC → instruction memory, fetch instruction
- Register numbers → register file, read registers
- Depending on instruction class
 - Use ALU to calculate
 - Arithmetic result
 - Memory address for load/store
 - Branch target address
 - Access data memory for load/store
 - PC ← target address or PC + 4

CPU Overview

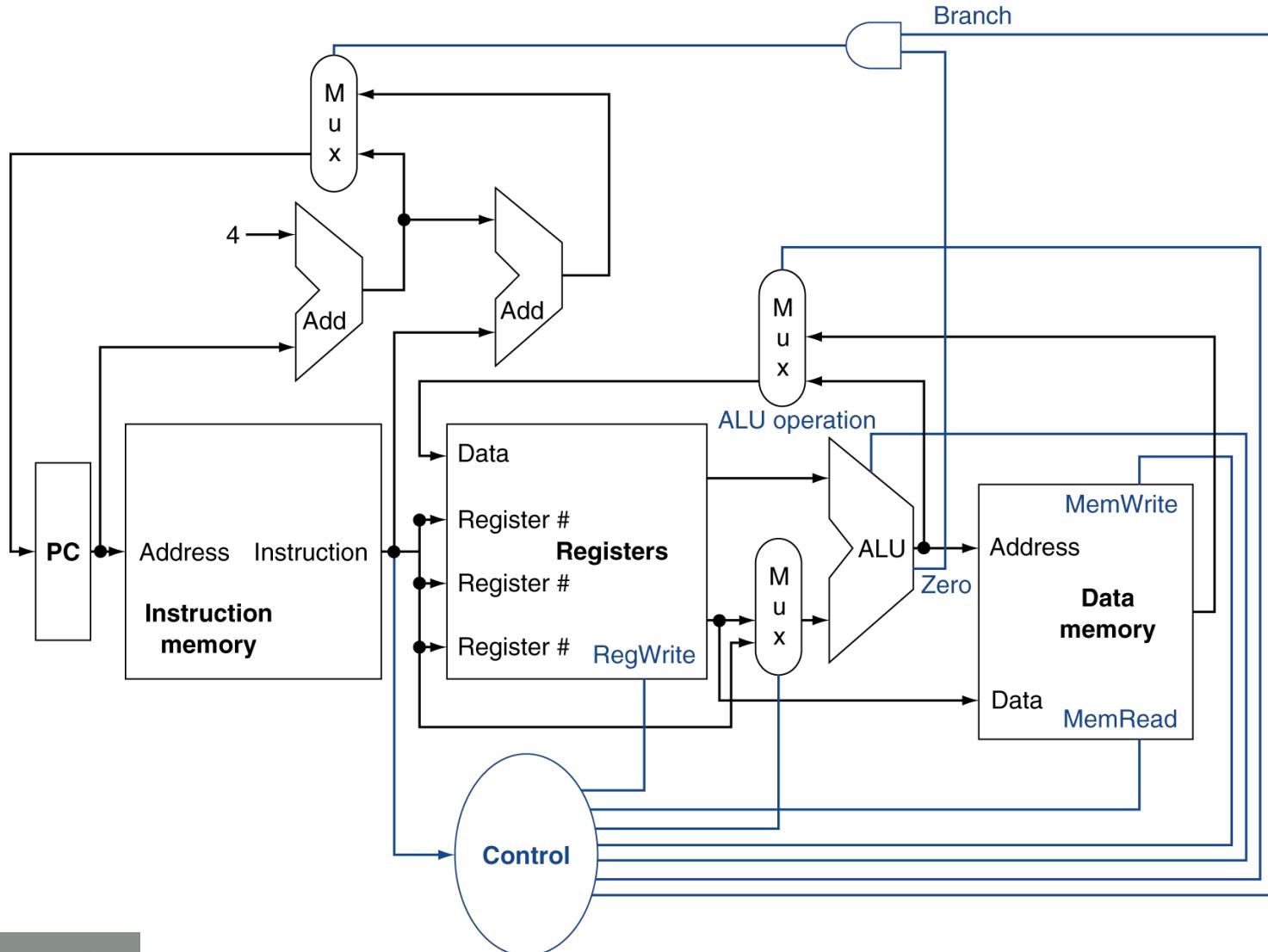


Multiplexers



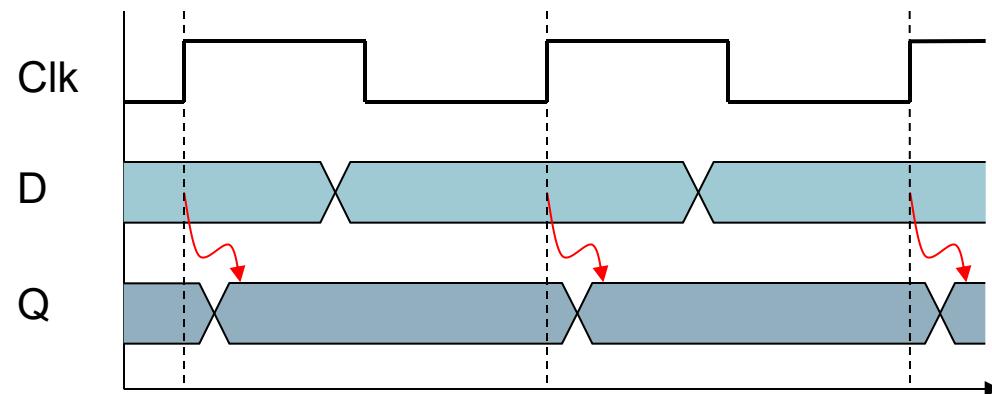
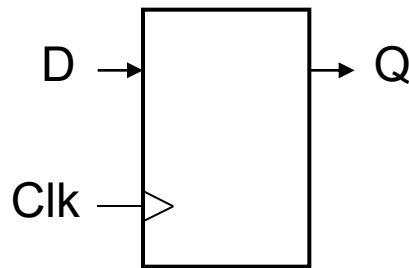
- Can't just join wires together
 - Use multiplexers

Control



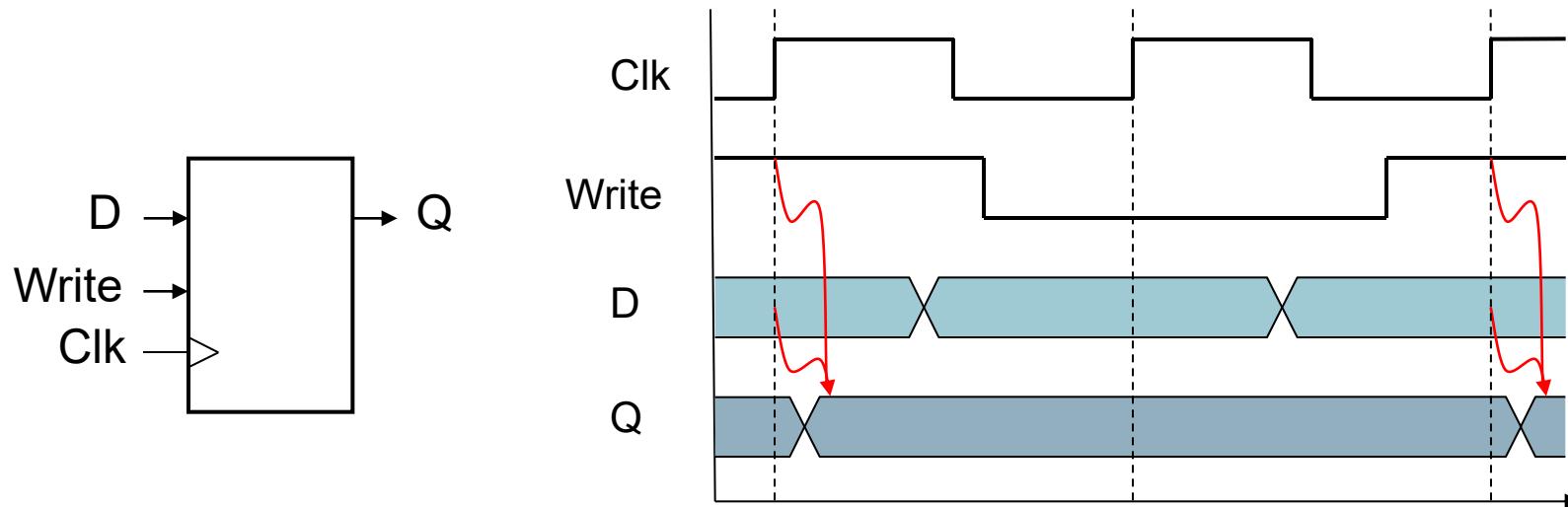
Sequential Elements

- Register: stores data in a circuit
 - Uses a clock signal to determine when to update the stored value
 - Edge-triggered: update when Clk changes from 0 to 1



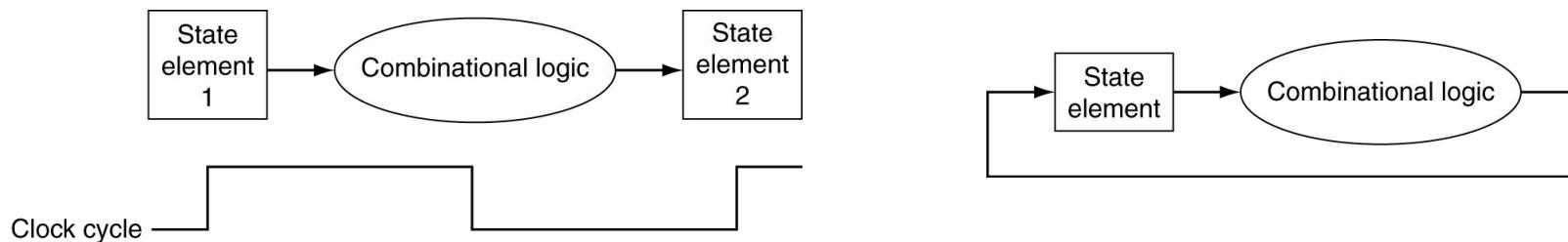
Sequential Elements

- Register with write control
 - Only updates on clock edge when write control input is 1
 - Used when stored value is required later



Clocking Methodology

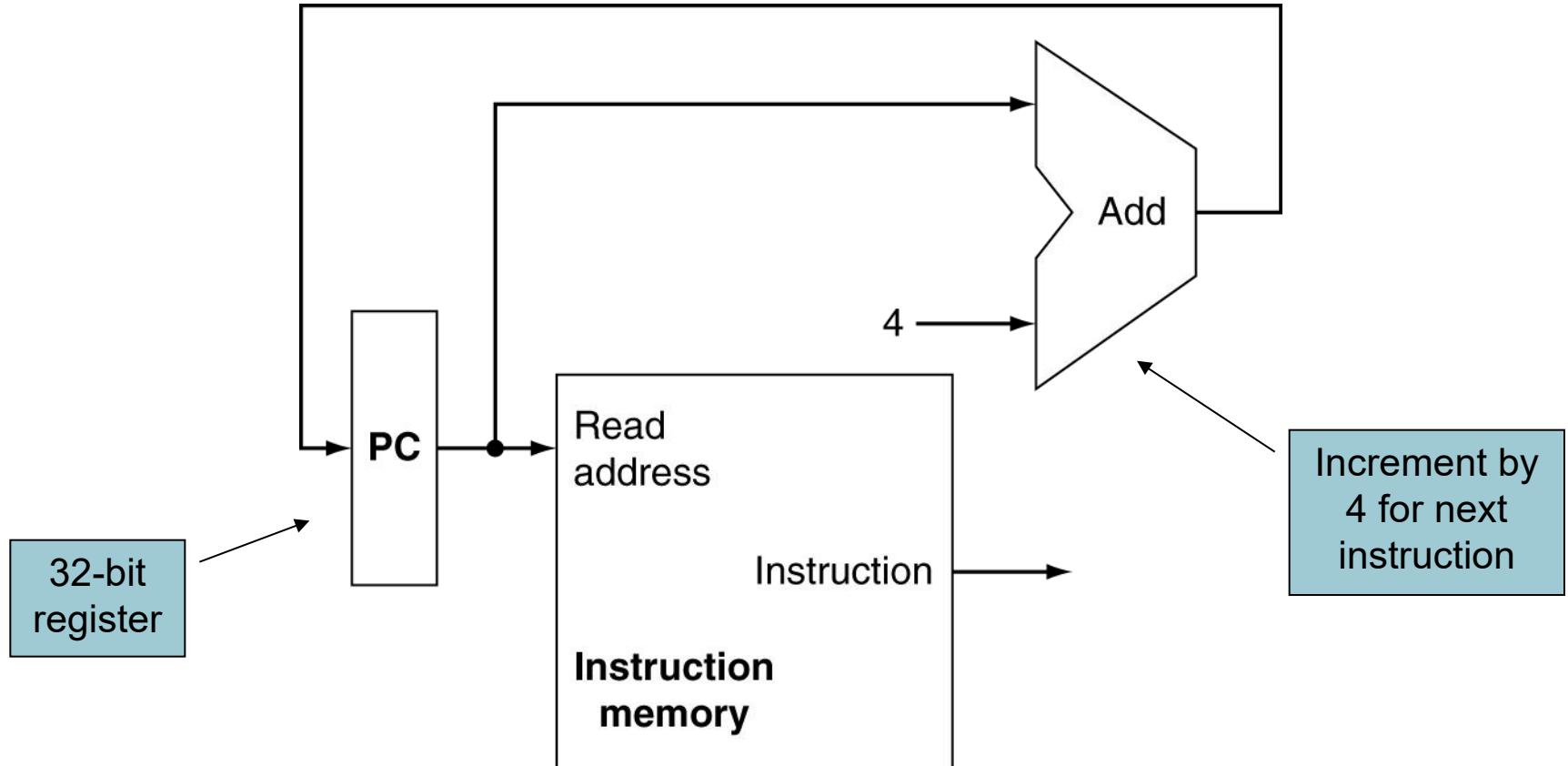
- Combinational logic transforms data during clock cycles
 - Between clock edges
 - Input from state elements, output to state element
 - Longest delay determines clock period



Building a Datapath

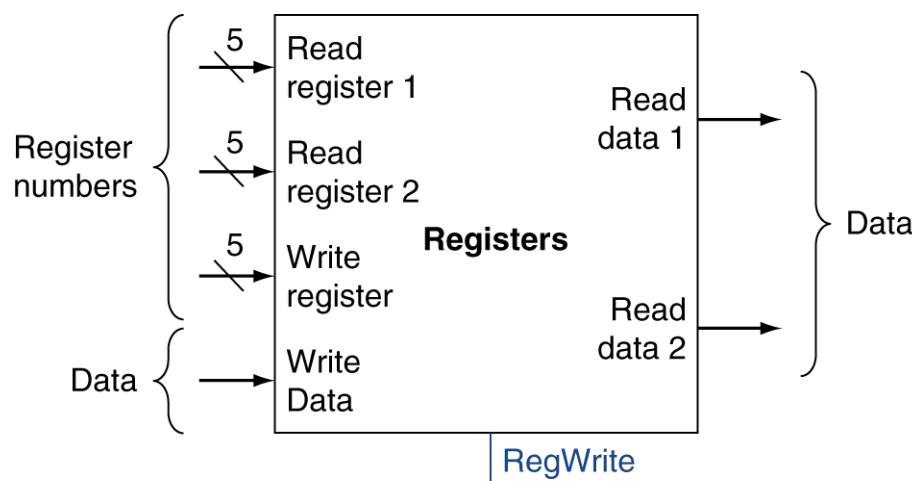
- Datapath
 - Elements that process data and addresses in the CPU
 - Registers, ALUs, mux's, memories, ...
- We will build a MIPS datapath incrementally
 - Refining the overview design

Instruction Fetch

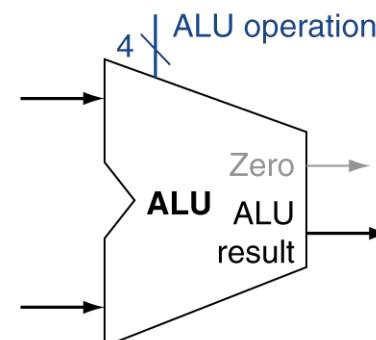


R-Format Instructions

- Read two register operands
- Perform arithmetic/logical operation
- Write register result



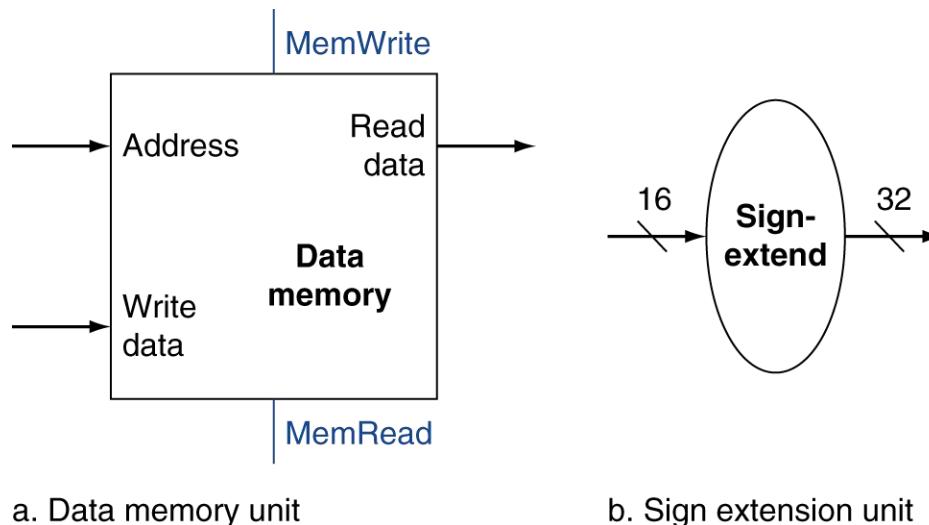
a. Registers



b. ALU

Load/Store Instructions

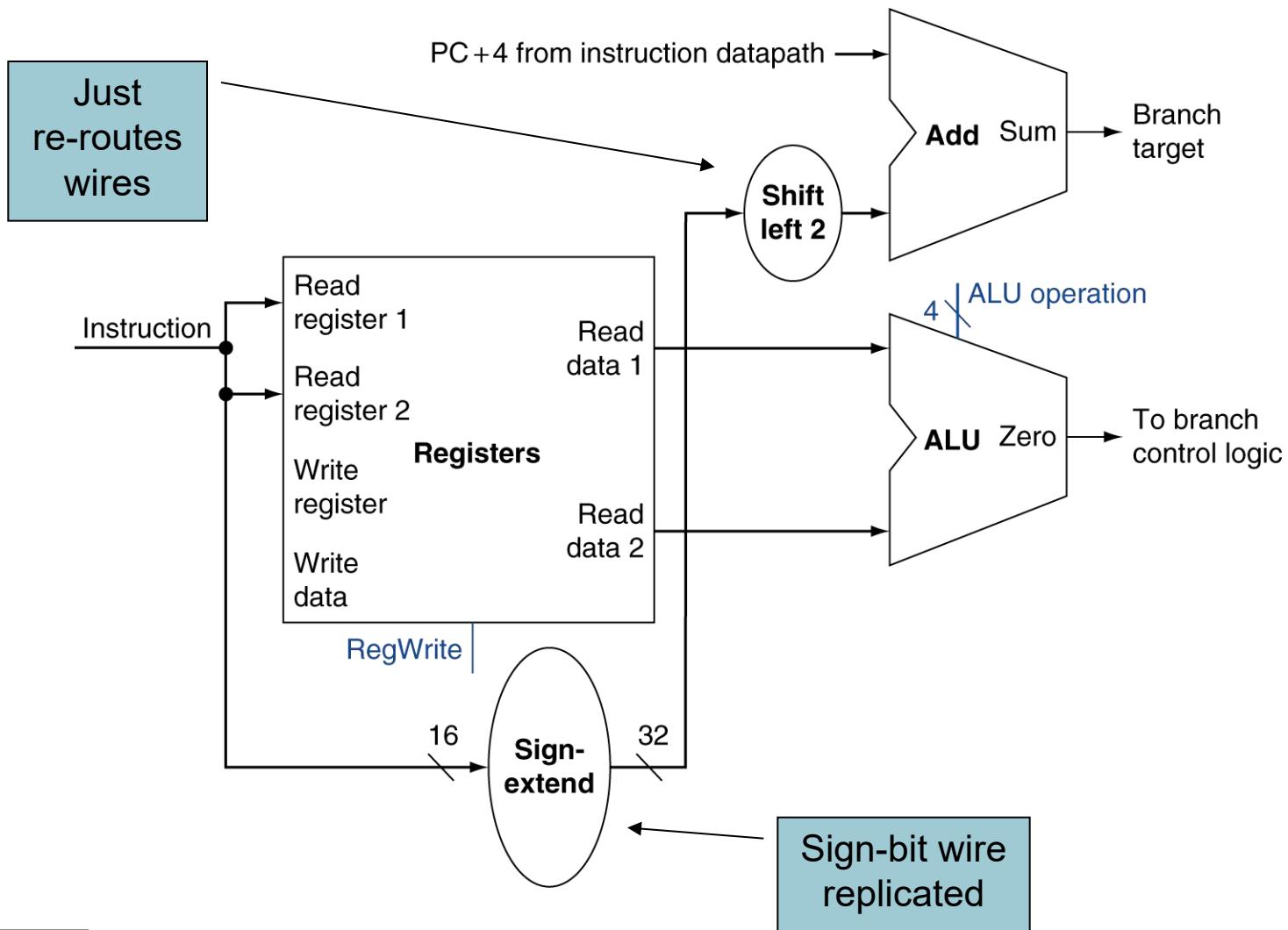
- Read register operands
- Calculate address using 16-bit offset
 - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory



Branch Instructions

- Read register operands
- Compare operands
 - Use ALU, subtract and check Zero output
- Calculate target address
 - Sign-extend displacement
 - Shift left 2 places (word displacement)
 - Add to PC + 4
 - Already calculated by instruction fetch

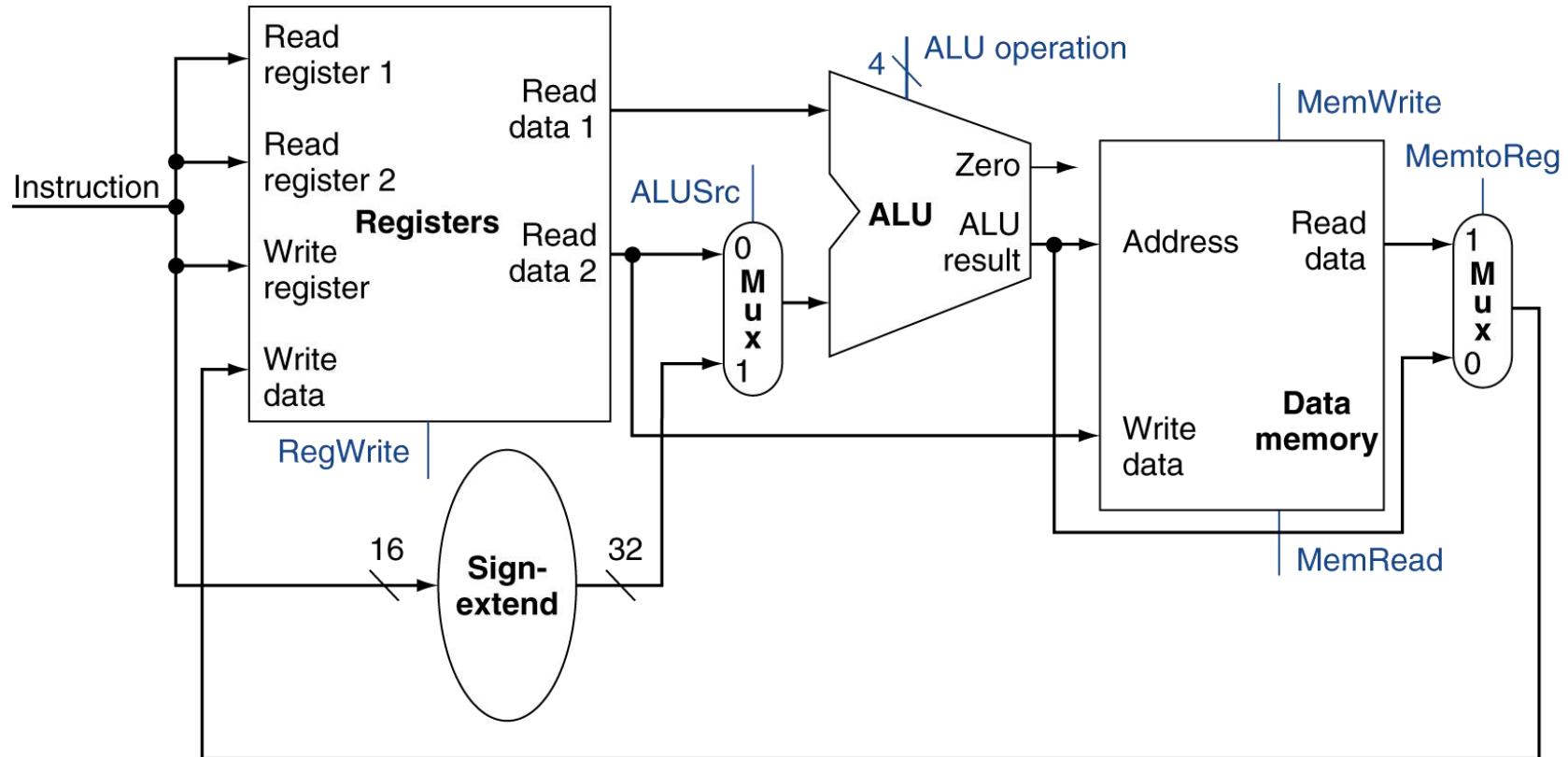
Branch Instructions



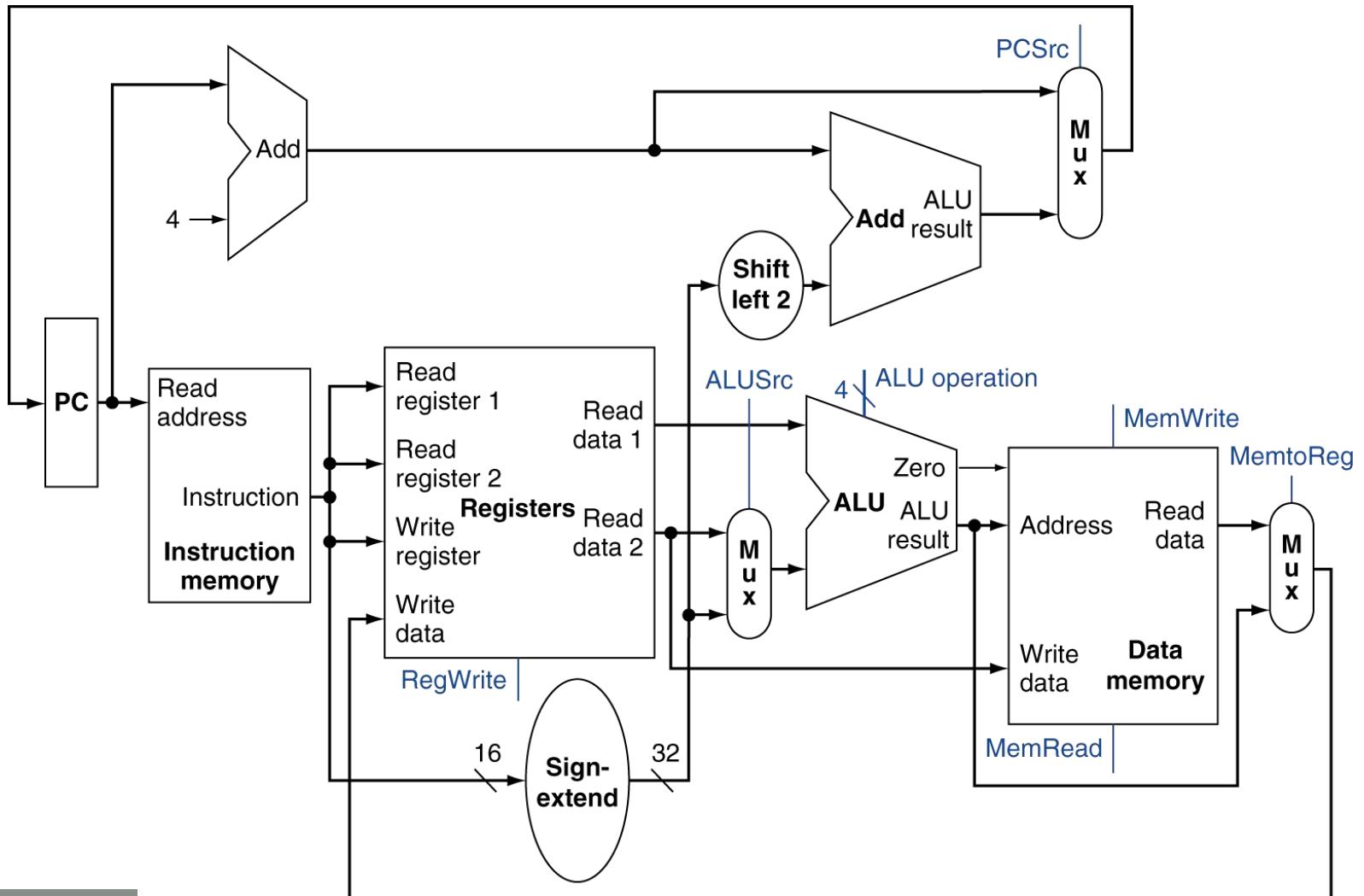
Composing the Elements

- First-cut data path does an instruction in one clock cycle
 - Each datapath element can only do one function at a time
 - Hence, we need separate instruction and data memories
- Use multiplexers where alternate data sources are used for different instructions

R-Type/Load/Store Datapath



Full Datapath



ALU Control

- ALU used for
 - Load/Store: $F = \text{add}$
 - Branch: $F = \text{subtract}$
 - R-type: F depends on funct field

ALU control	Function
0000	AND
0001	OR
0010	add
0110	subtract
0111	set-on-less-than
1100	NOR

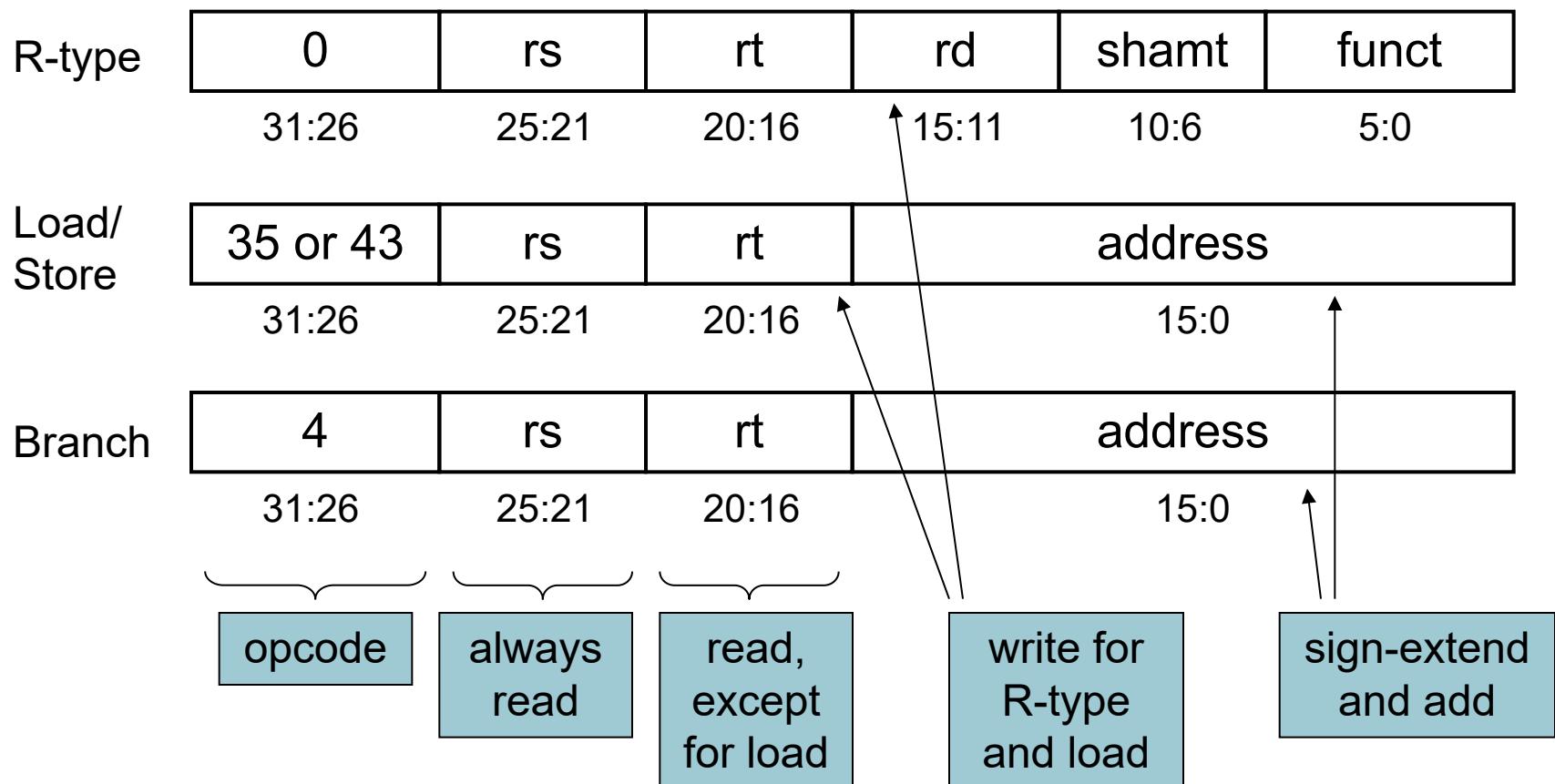
ALU Control

- Assume 2-bit ALUOp derived from opcode
 - Combinational logic derives ALU control

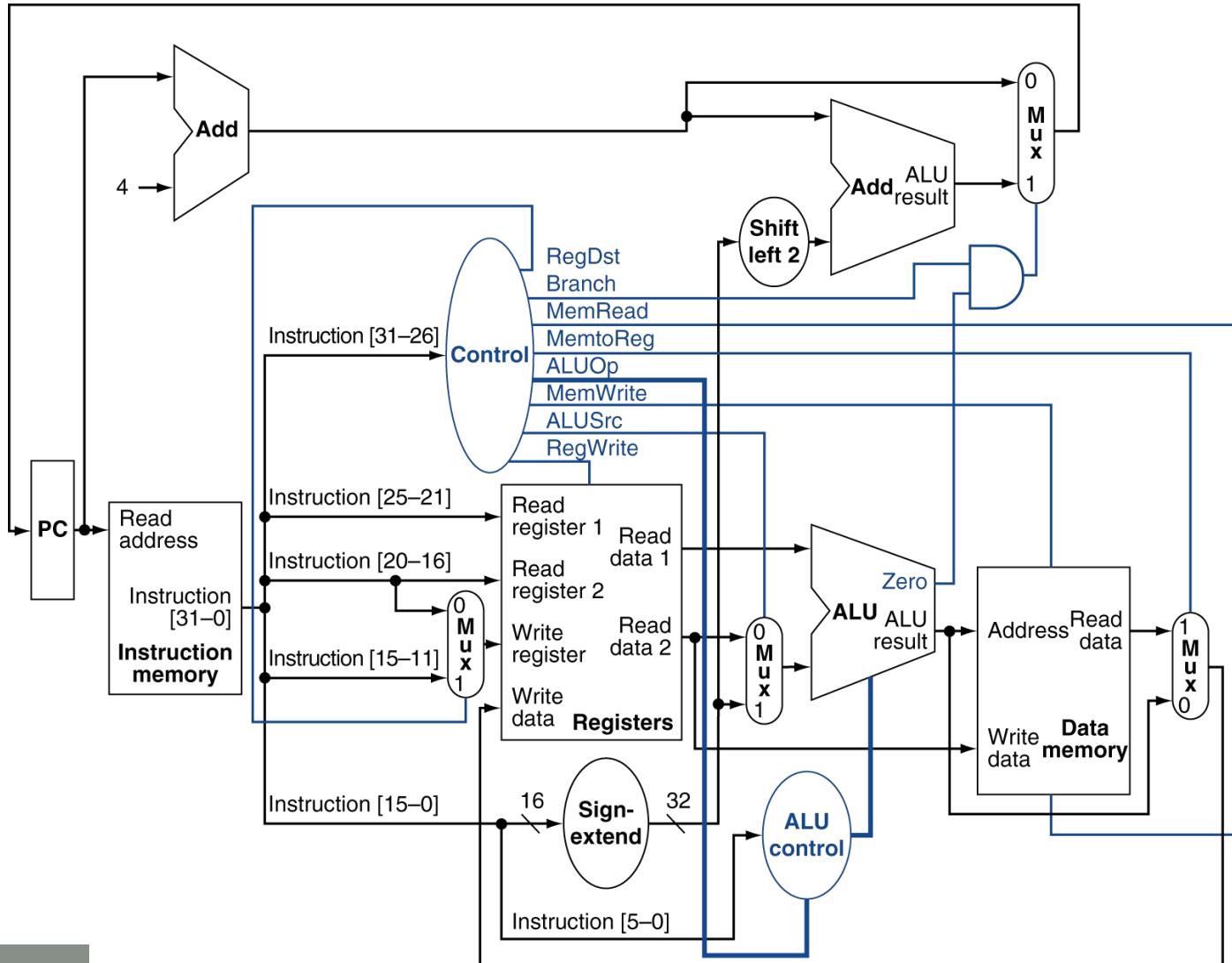
opcode	ALUOp	Operation	funct	ALU function	ALU control
lw	00	load word	XXXXXX	add	0010
sw	00	store word	XXXXXX	add	0010
beq	01	branch equal	XXXXXX	subtract	0110
R-type	10	add	100000	add	0010
		subtract	100010	subtract	0110
		AND	100100	AND	0000
		OR	100101	OR	0001
		set-on-less-than	101010	set-on-less-than	0111

The Main Control Unit

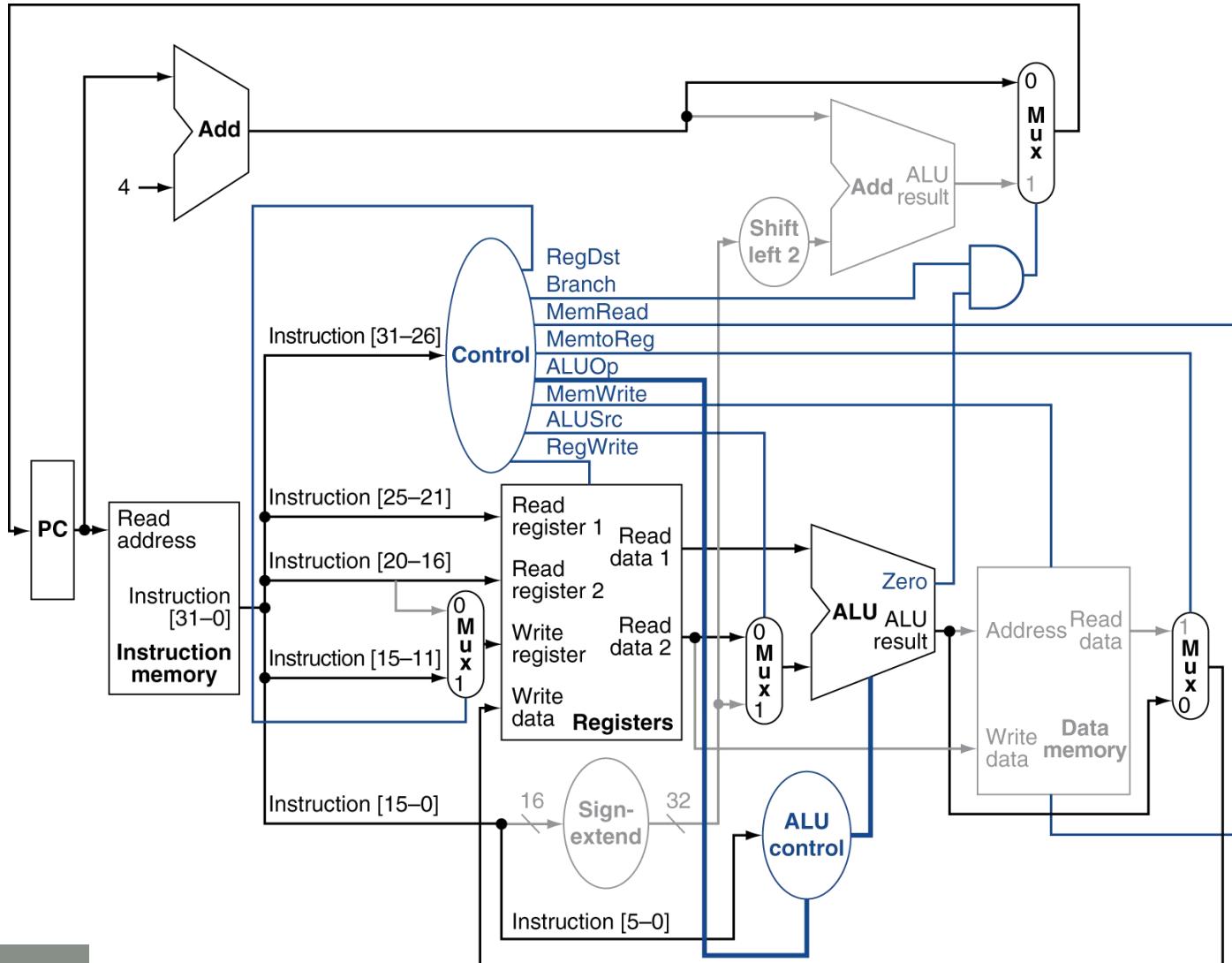
- Control signals derived from instruction



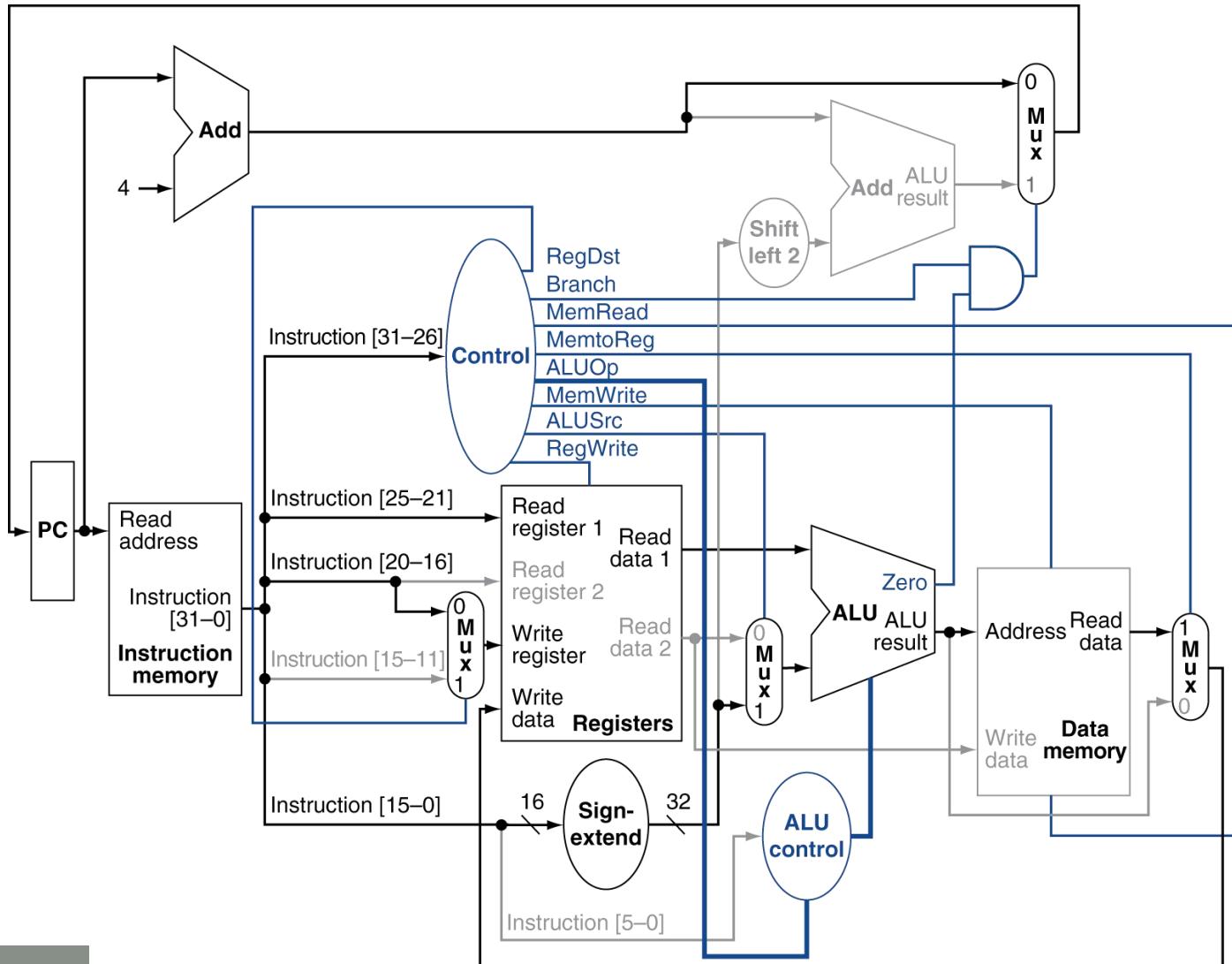
Datapath With Control



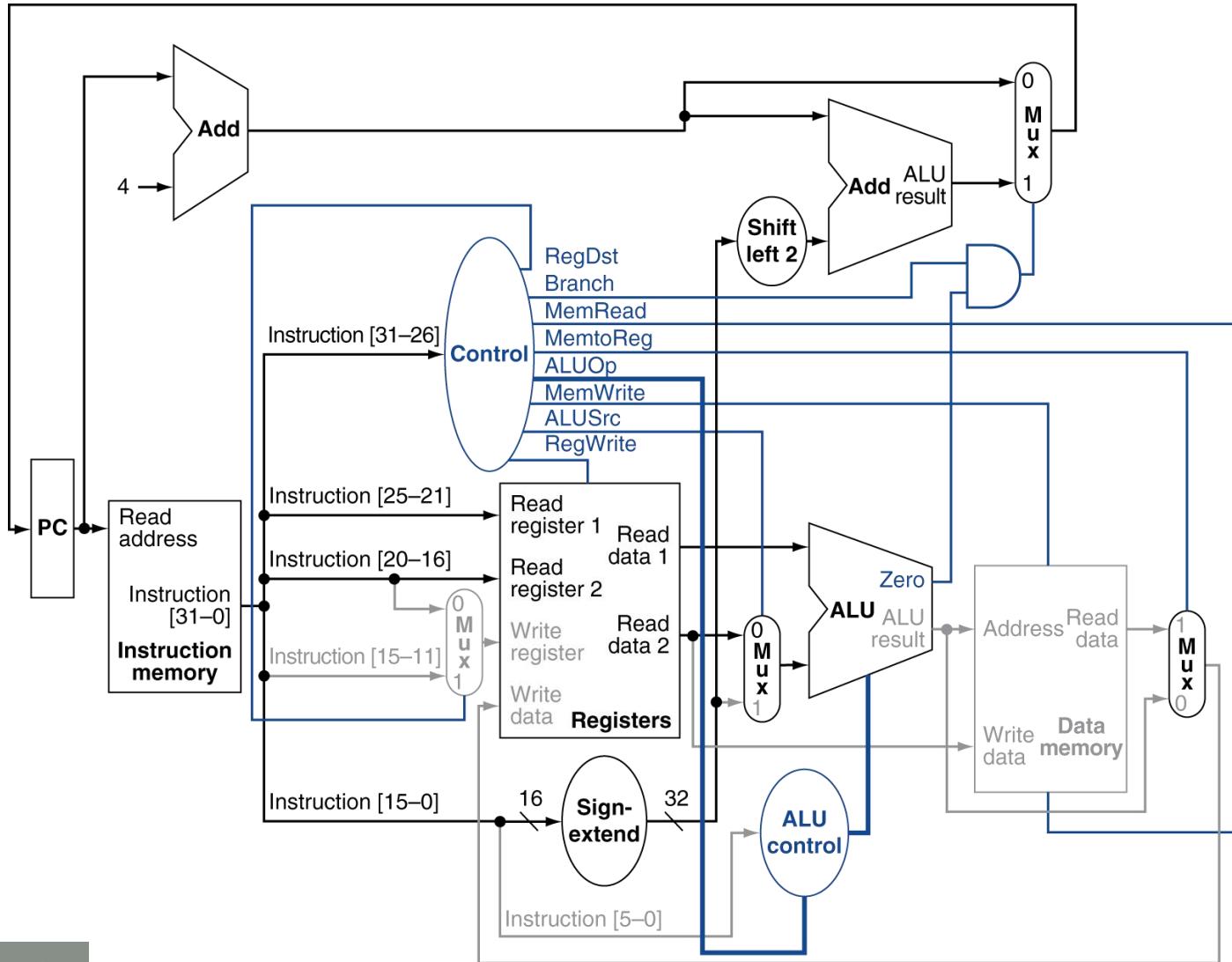
R-Type Instruction



Load Instruction

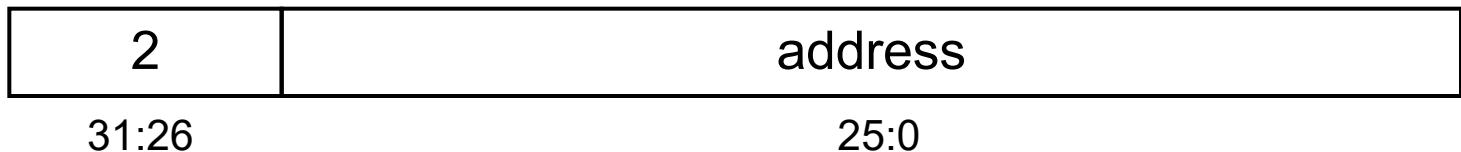


Branch-on-Equal Instruction



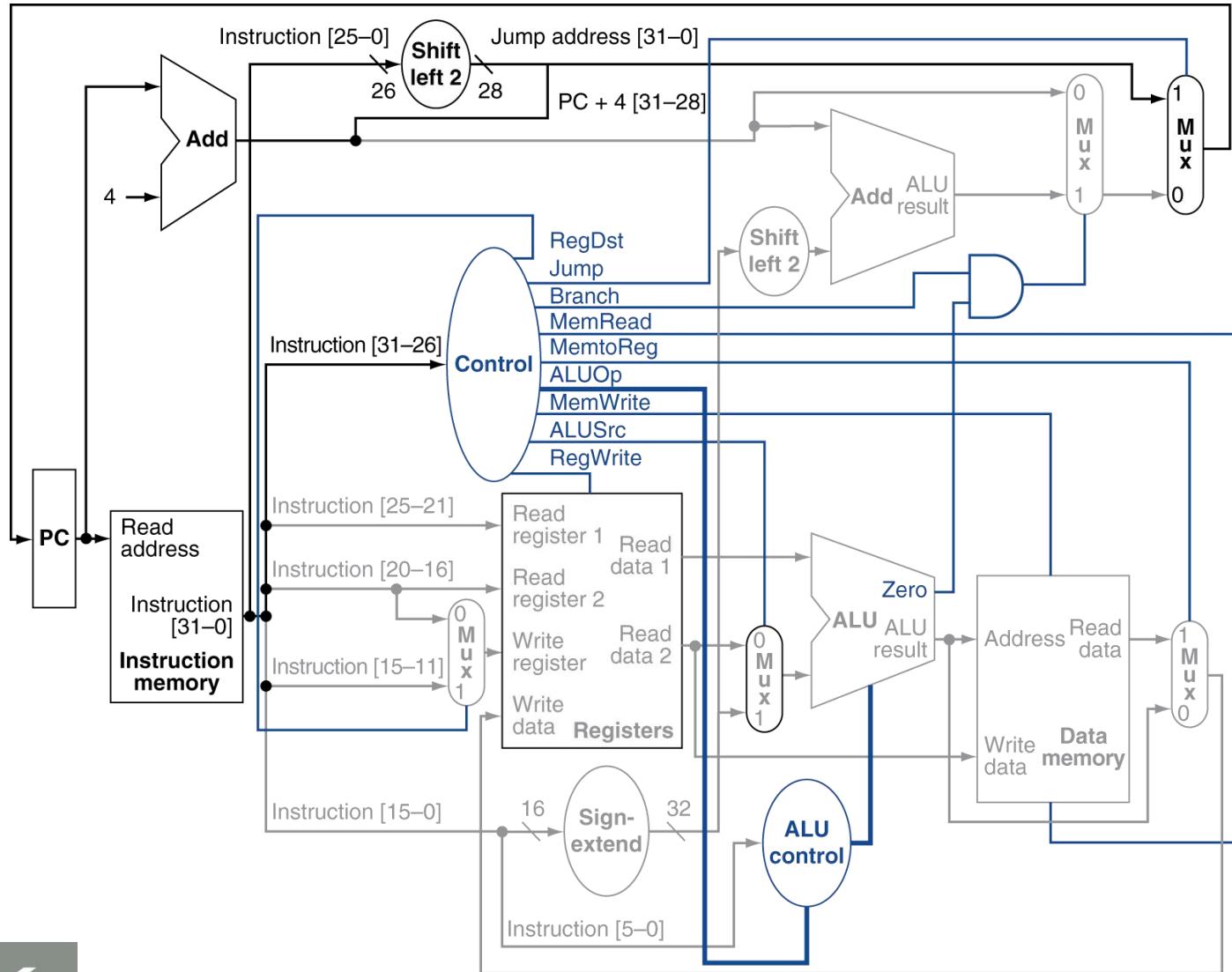
Implementing Jumps

Jump



- Jump uses word address
- Update PC with concatenation of
 - Top 4 bits of old PC
 - 26-bit jump address
 - 00
- Need an extra control signal decoded from opcode

Datapath With Jumps Added

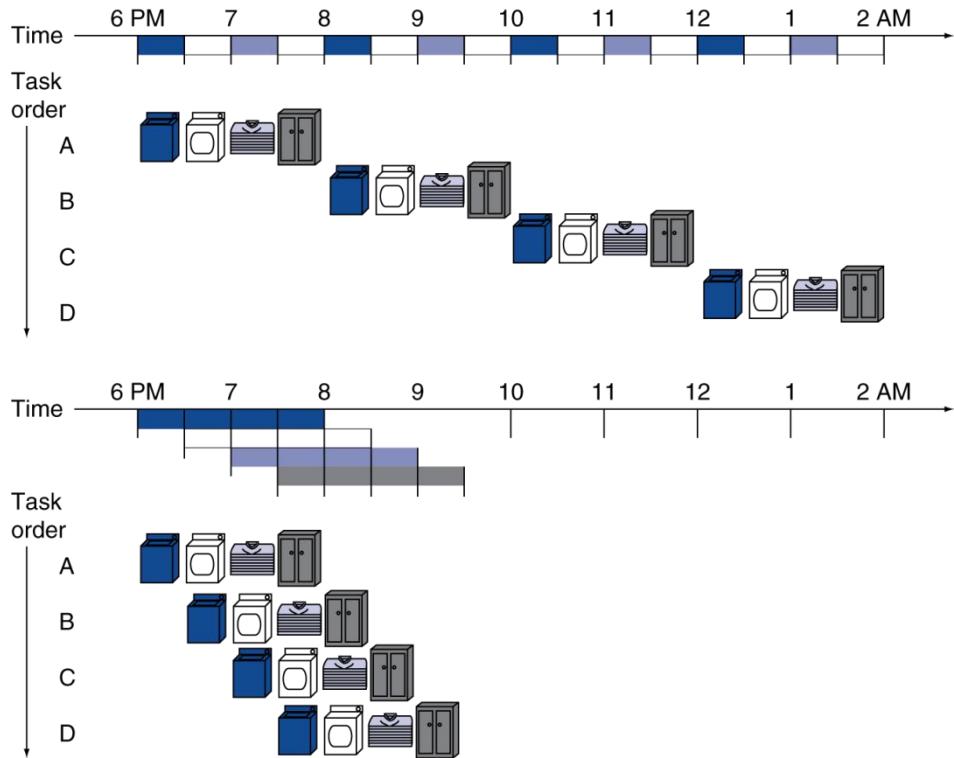


Performance Issues

- Longest delay determines clock period
 - Critical path: load instruction
 - Instruction memory → register file → ALU → data memory → register file
- Not feasible to vary period for different instructions
- Violates design principle
 - Making the common case fast
- We will improve performance by pipelining

Pipelining Analogy

- Pipelined laundry: overlapping execution
 - Parallelism improves performance



- Four loads:
 - Speedup
 $= 8/3.5 = 2.3$
- Non-stop:
 - Speedup
 $= 2n/0.5n + 1.5 \approx 4$
= number of stages

MIPS Pipeline

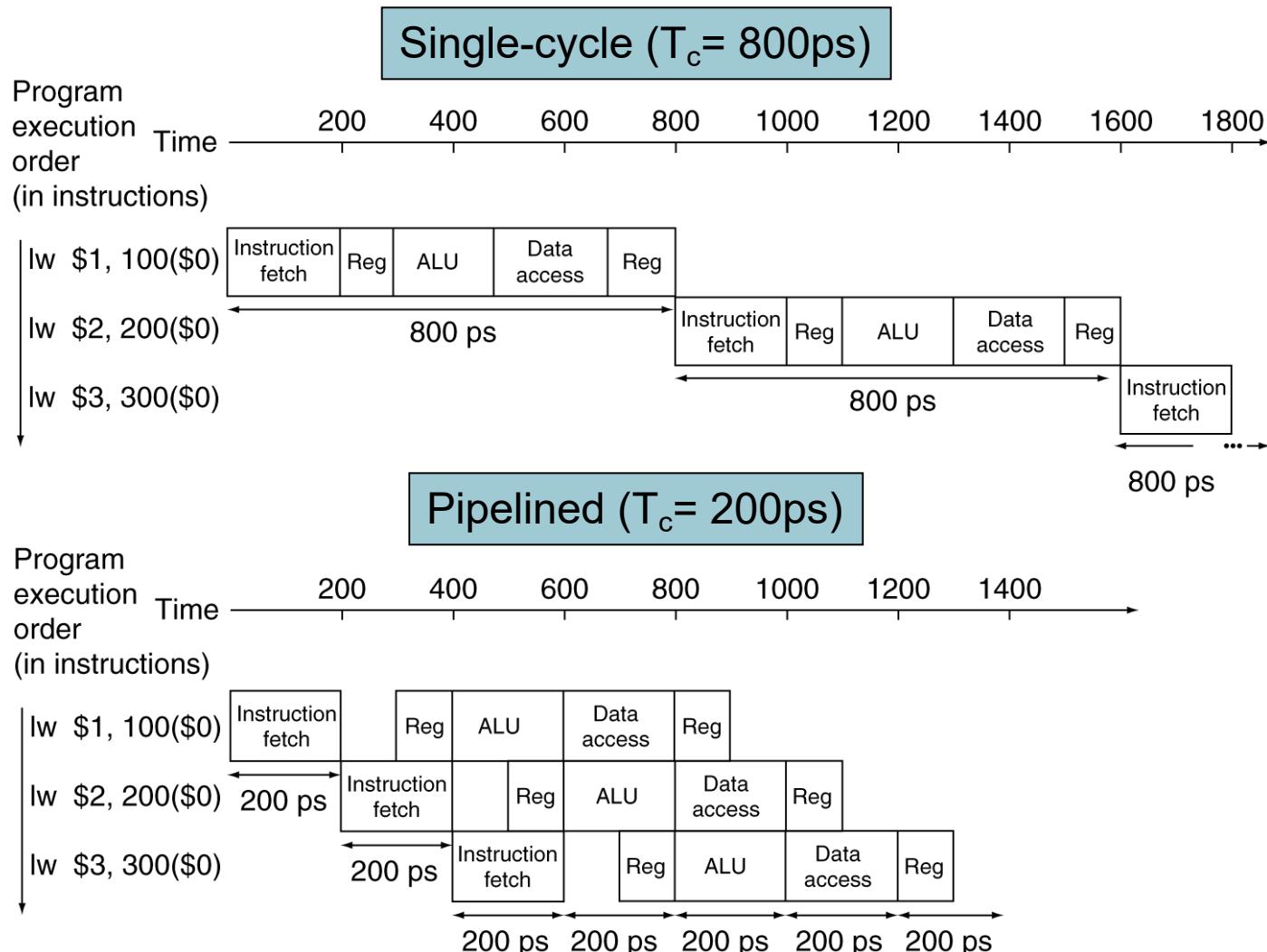
- Five stages, one step per stage
 1. IF: Instruction fetch from memory
 2. ID: Instruction decode & register read
 3. EX: Execute operation or calculate address
 4. MEM: Access memory operand
 5. WB: Write result back to register

Pipeline Performance

- Assume time for stages is
 - 100ps for register read or write
 - 200ps for other stages
- Compare pipelined datapath with single-cycle datapath

Instr	Instr fetch	Register read	ALU op	Memory access	Register write	Total time
lw	200ps	100 ps	200ps	200ps	100 ps	800ps
sw	200ps	100 ps	200ps	200ps		700ps
R-format	200ps	100 ps	200ps		100 ps	600ps
beq	200ps	100 ps	200ps			500ps

Pipeline Performance



Pipeline Speedup

- If all stages are balanced
 - i.e., all take the same time
 - Time between instructions_{pipelined}
= Time between instructions_{nonpipelined}

Number of stages
- If not balanced, speedup is less
- Speedup due to increased throughput
 - Latency (time for each instruction) does not decrease

Pipelining and ISA Design

- MIPS ISA designed for pipelining
 - All instructions are 32-bits
 - Easier to fetch and decode in one cycle
 - c.f. x86: 1- to 17-byte instructions
 - Few and regular instruction formats
 - Can decode and read registers in one step
 - Load/store addressing
 - Can calculate address in 3rd stage, access memory in 4th stage
 - Alignment of memory operands
 - Memory access takes only one cycle

Hazards

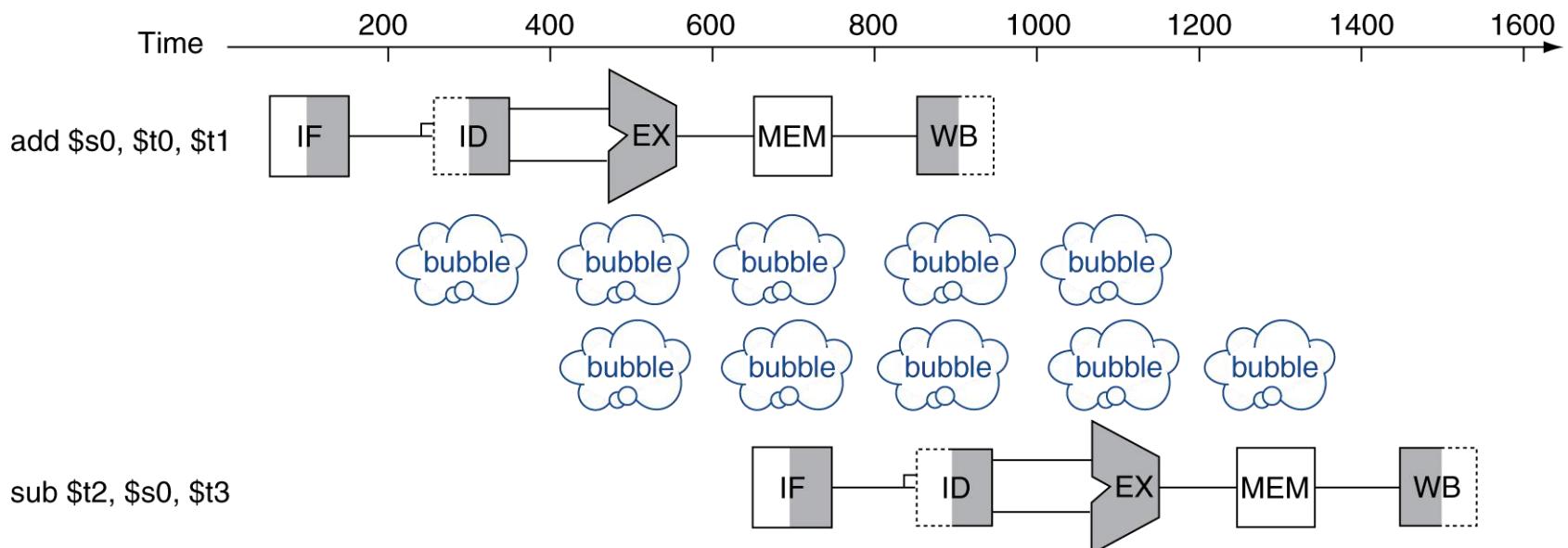
- Situations that prevent starting the next instruction in the next cycle
- Structure hazards
 - A required resource is busy
- Data hazard
 - Need to wait for previous instruction to complete its data read/write
- Control hazard
 - Deciding on control action depends on previous instruction

Structure Hazards

- Conflict for use of a resource
- In MIPS pipeline with a single memory
 - Load/store requires data access
 - Instruction fetch would have to *stall* for that cycle
 - Would cause a pipeline “bubble”
- Hence, pipelined datapaths require separate instruction/data memories
 - Or separate instruction/data caches

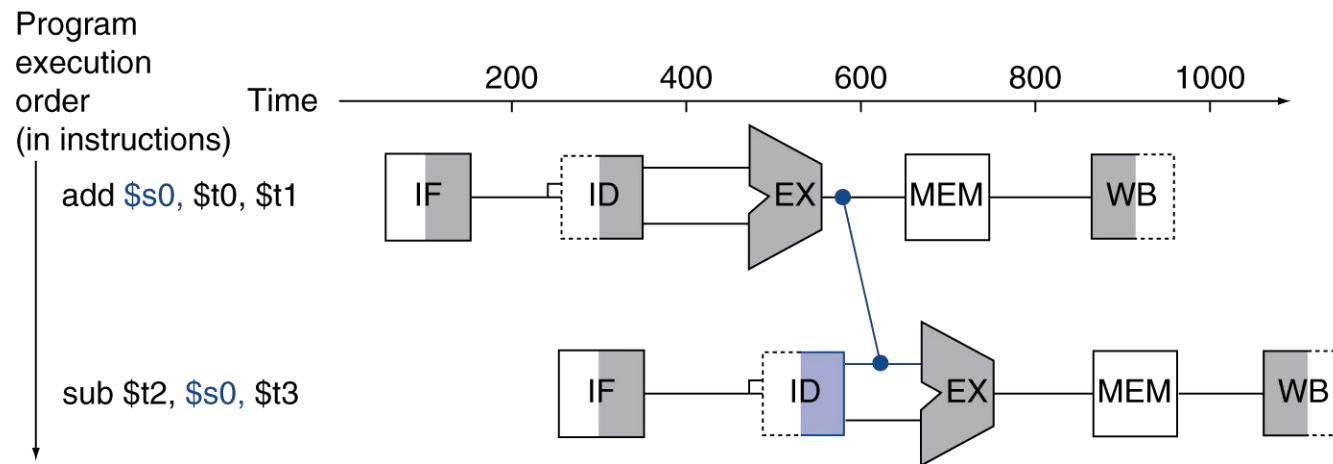
Data Hazards

- An instruction depends on completion of data access by a previous instruction
 - add \$s0, \$t0, \$t1
 - sub \$t2, \$s0, \$t3



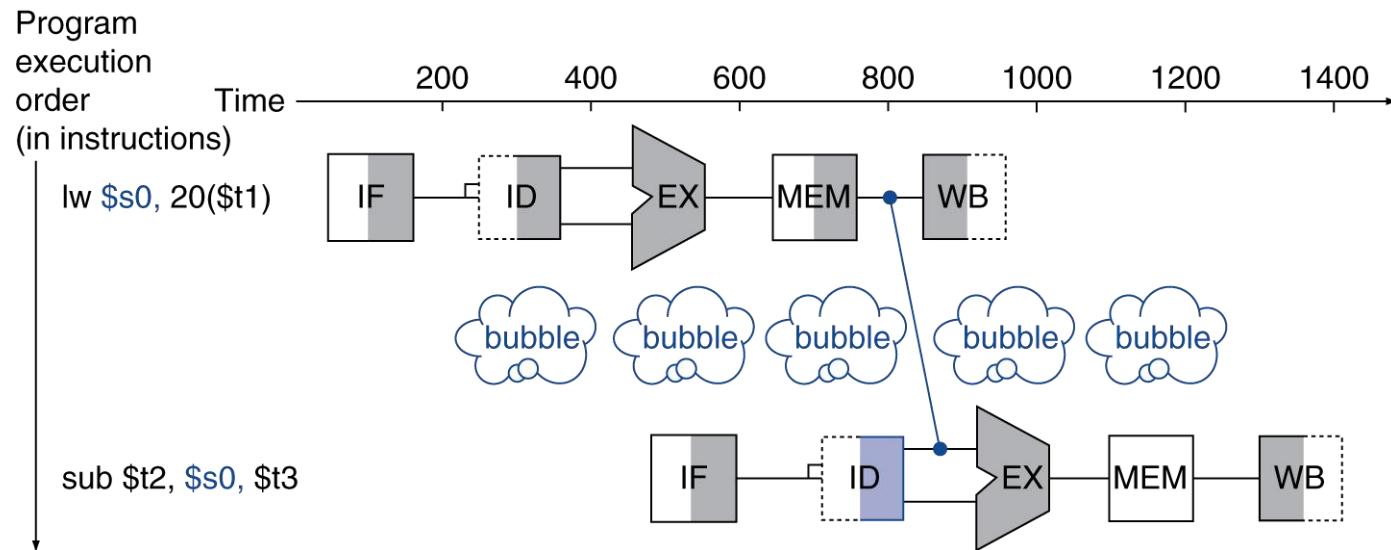
Forwarding (aka Bypassing)

- Use result when it is computed
 - Don't wait for it to be stored in a register
 - Requires extra connections in the datapath



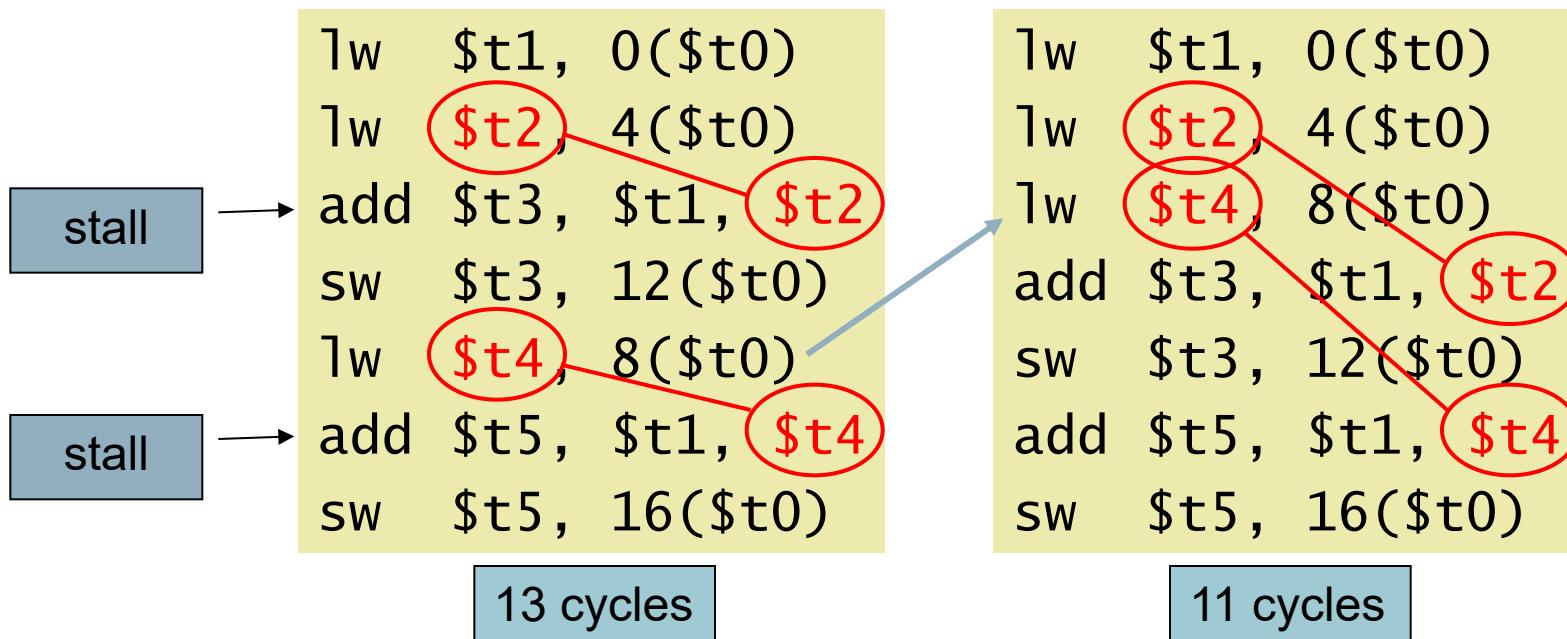
Load-Use Data Hazard

- Can't always avoid stalls by forwarding
 - If value not computed when needed
 - Can't forward backward in time!



Code Scheduling to Avoid Stalls

- Reorder code to avoid use of load result in the next instruction
- C code for $A = B + E; C = B + F;$

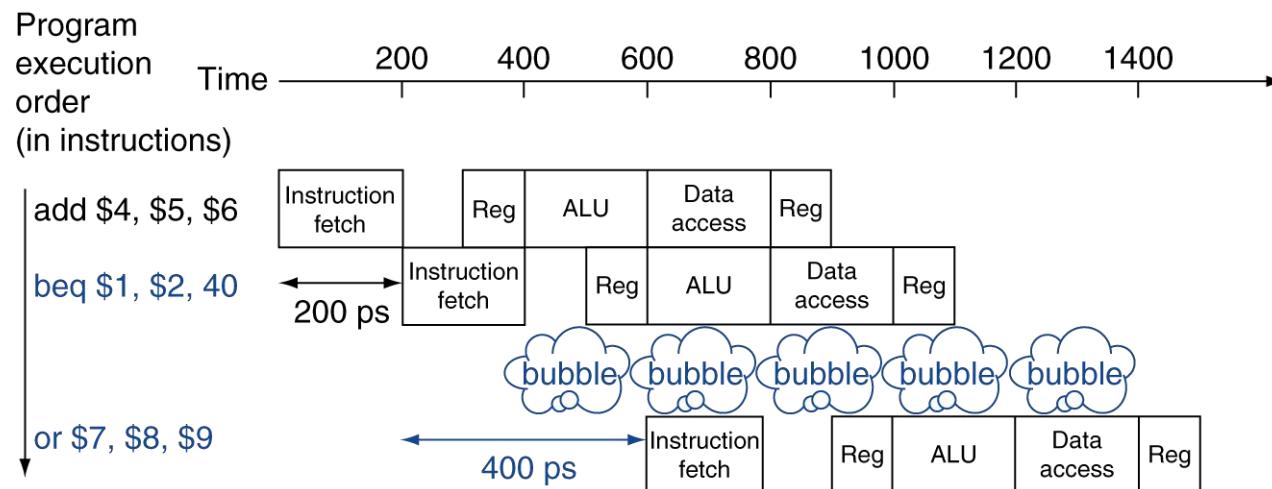


Control Hazards

- Branch determines flow of control
 - Fetching next instruction depends on branch outcome
 - Pipeline can't always fetch correct instruction
 - Still working on ID stage of branch
- In MIPS pipeline
 - Need to compare registers and compute target early in the pipeline
 - Add hardware to do it in ID stage

Stall on Branch

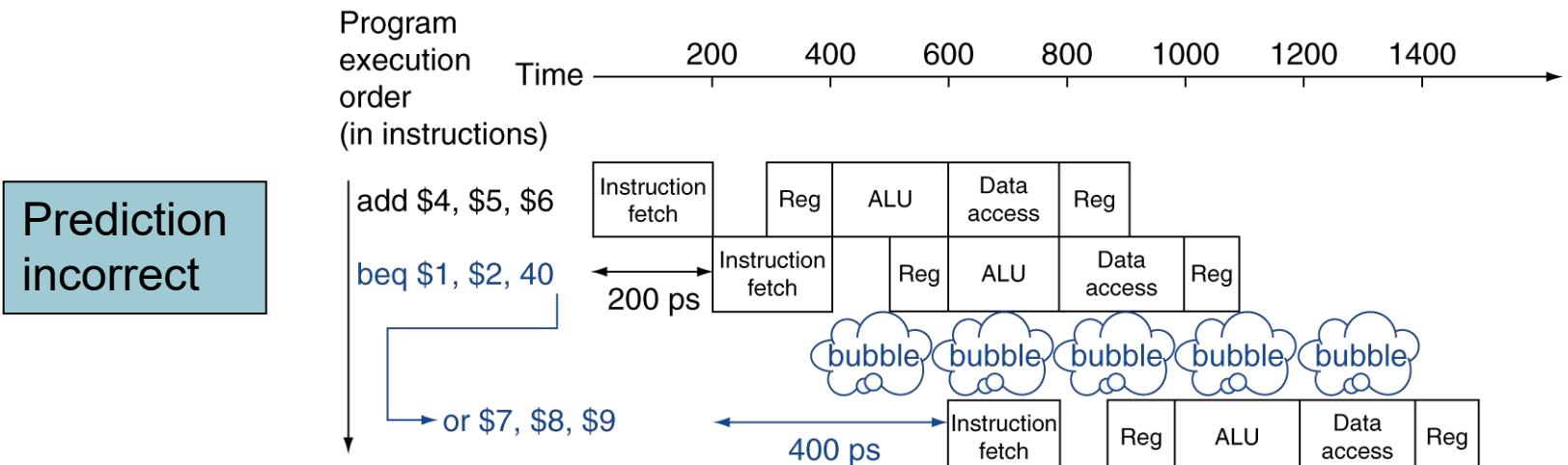
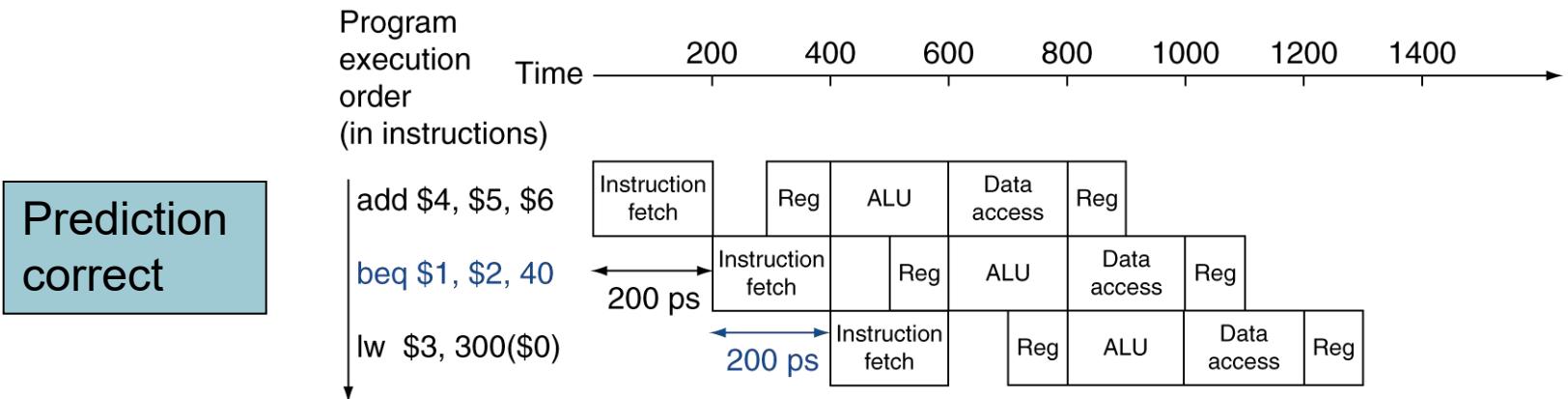
- Wait until branch outcome determined before fetching next instruction



Branch Prediction

- Longer pipelines can't readily determine branch outcome early
 - Stall penalty becomes unacceptable
- Predict outcome of branch
 - Only stall if prediction is wrong
- In MIPS pipeline
 - Can predict branches not taken
 - Fetch instruction after branch, with no delay

MIPS with Predict Not Taken



More-Realistic Branch Prediction

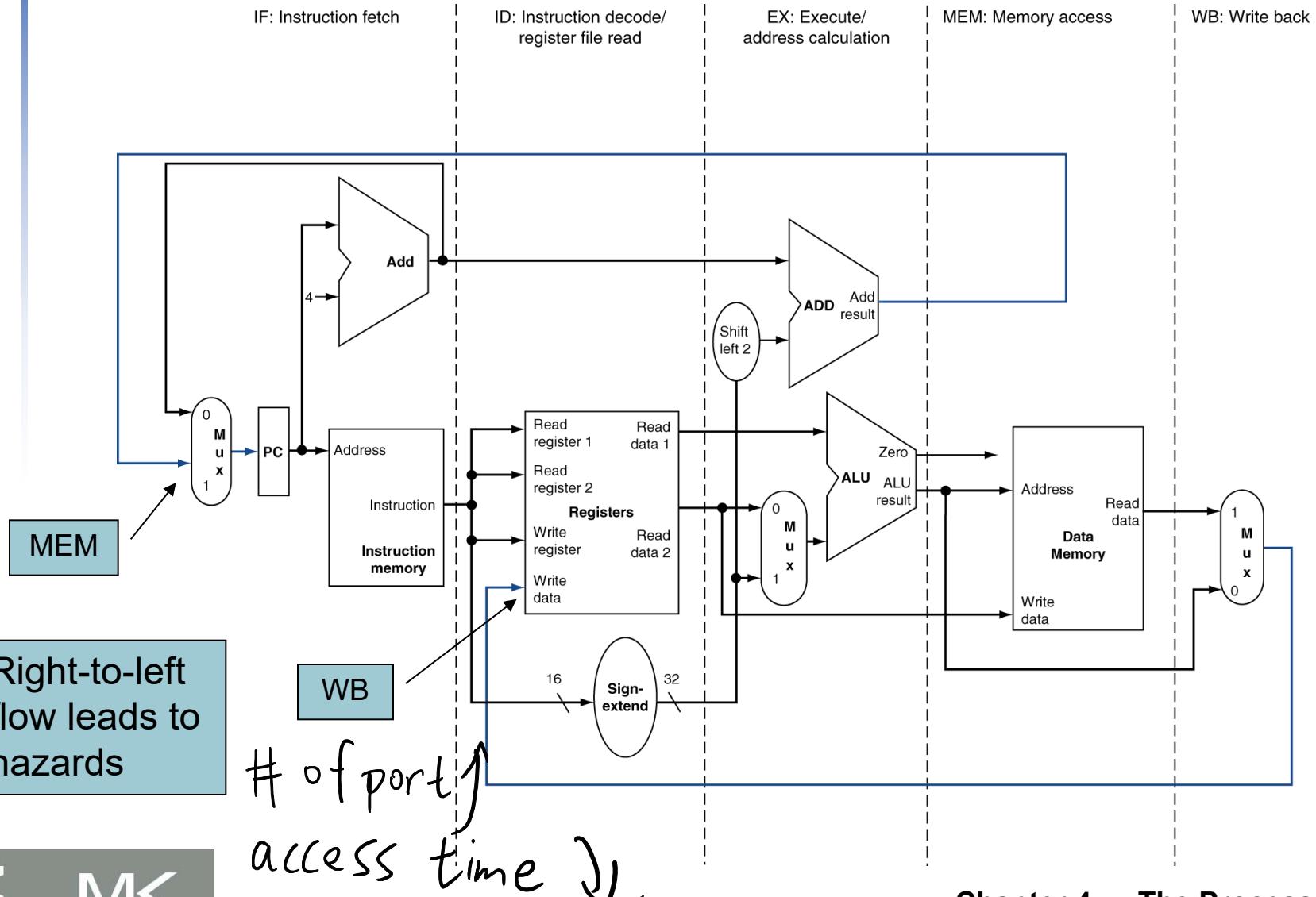
- Static branch prediction
 - Based on typical branch behavior
 - Example: loop and if-statement branches
 - Predict backward branches taken
 - Predict forward branches not taken
- Dynamic branch prediction
 - Hardware measures actual branch behavior
 - e.g., record recent history of each branch
 - Assume future behavior will continue the trend
 - When wrong, stall while re-fetching, and update history

Pipeline Summary

The BIG Picture

- Pipelining improves performance by increasing instruction throughput
 - Executes multiple instructions in parallel
 - Each instruction has the same latency
- Subject to hazards
 - Structure, data, control
- Instruction set design affects complexity of pipeline implementation

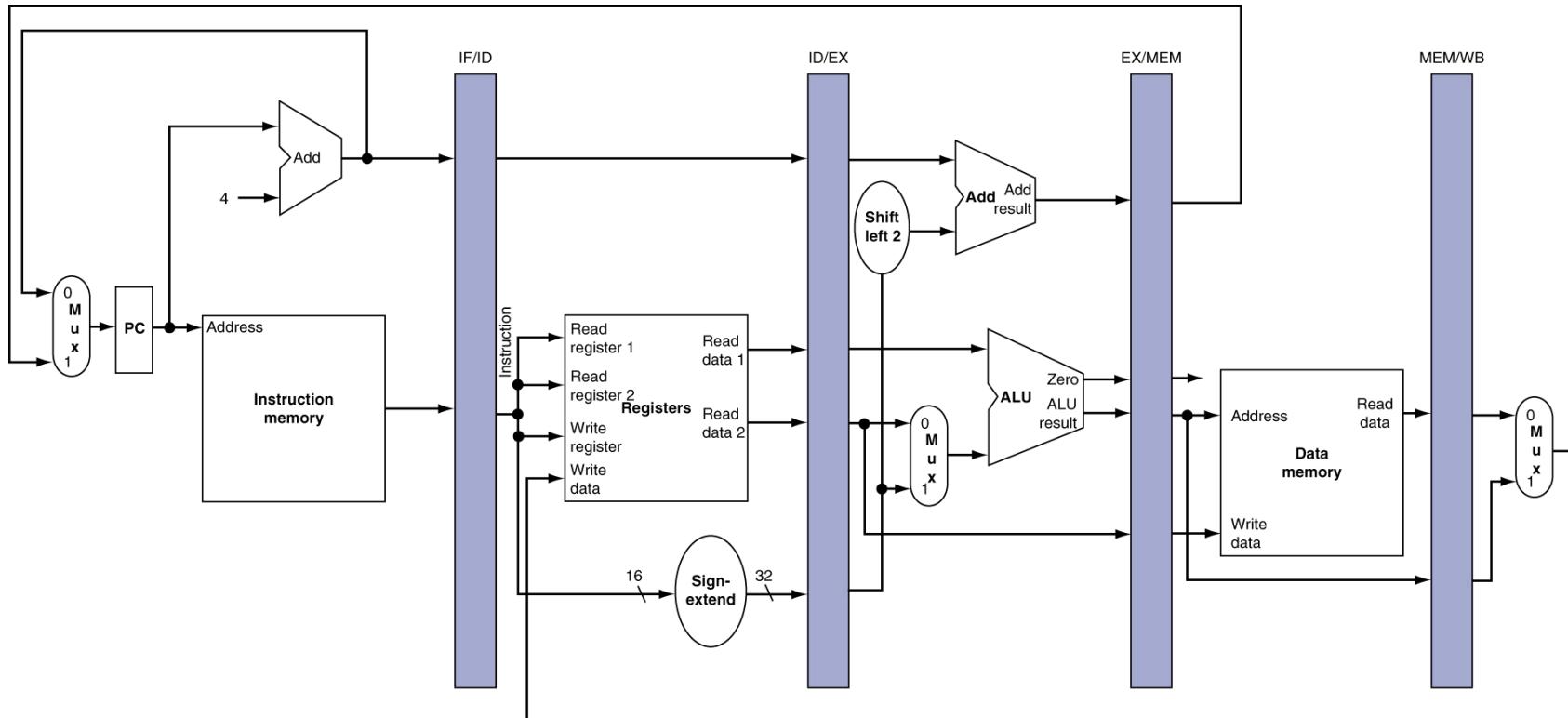
MIPS Pipelined Datapath



Pipeline registers

Intel 64 bit Data
64 bit instruction

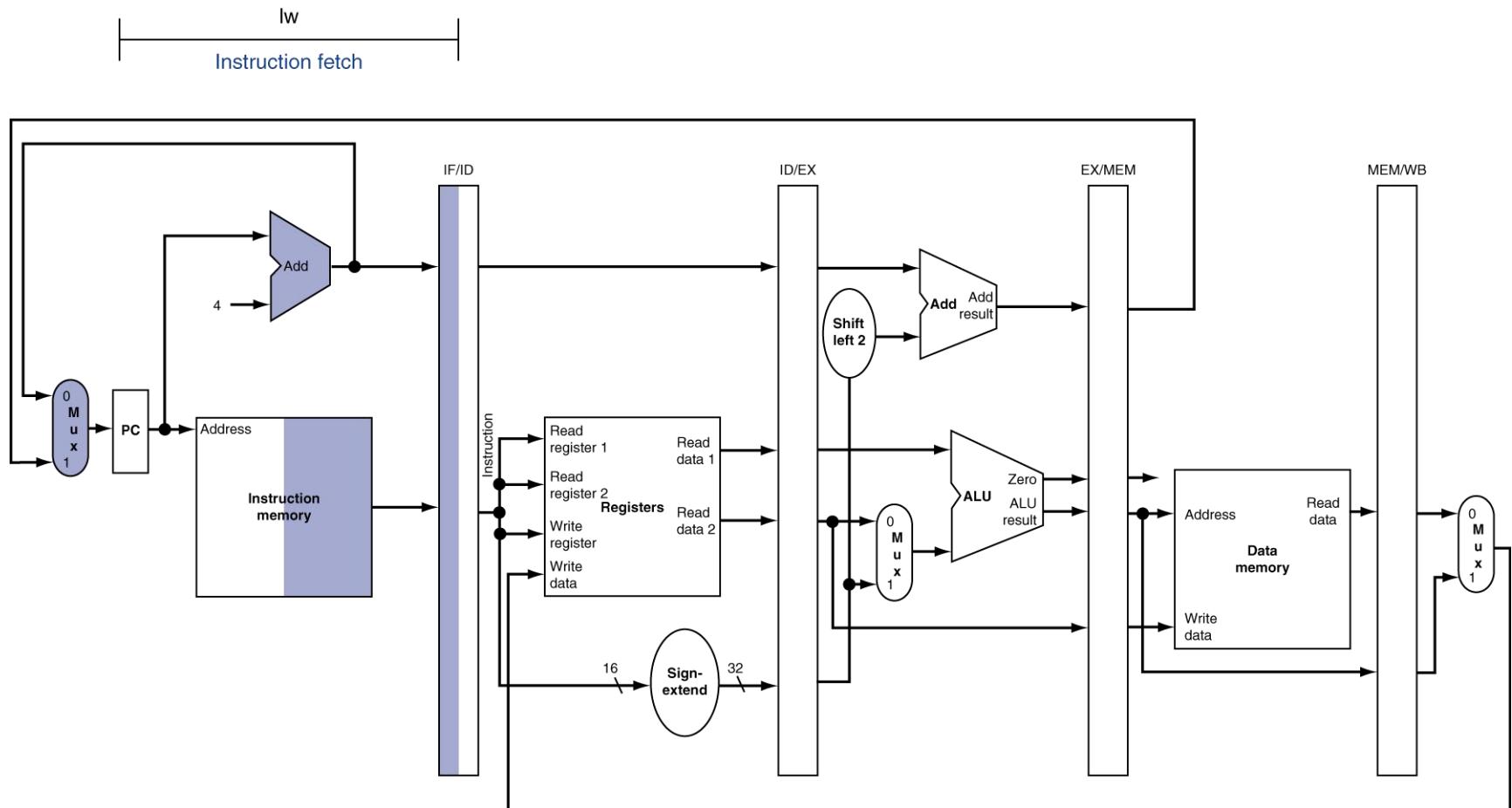
- Need registers between stages
 - To hold information produced in previous cycle



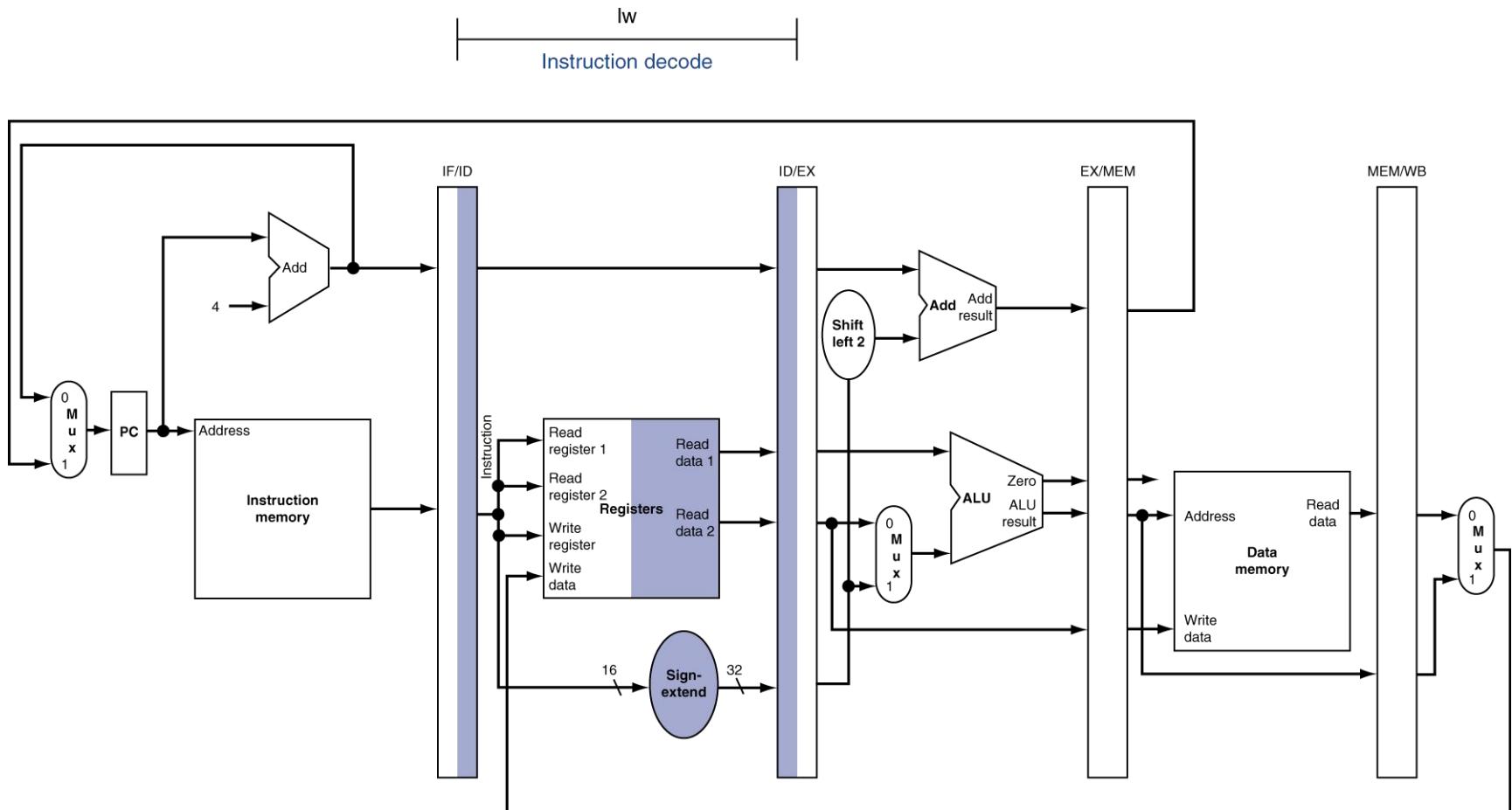
Pipeline Operation

- Cycle-by-cycle flow of instructions through the pipelined datapath
 - “Single-clock-cycle” pipeline diagram
 - Shows pipeline usage in a single cycle
 - Highlight resources used
 - c.f. “multi-clock-cycle” diagram
 - Graph of operation over time
- We’ll look at “single-clock-cycle” diagrams for load & store

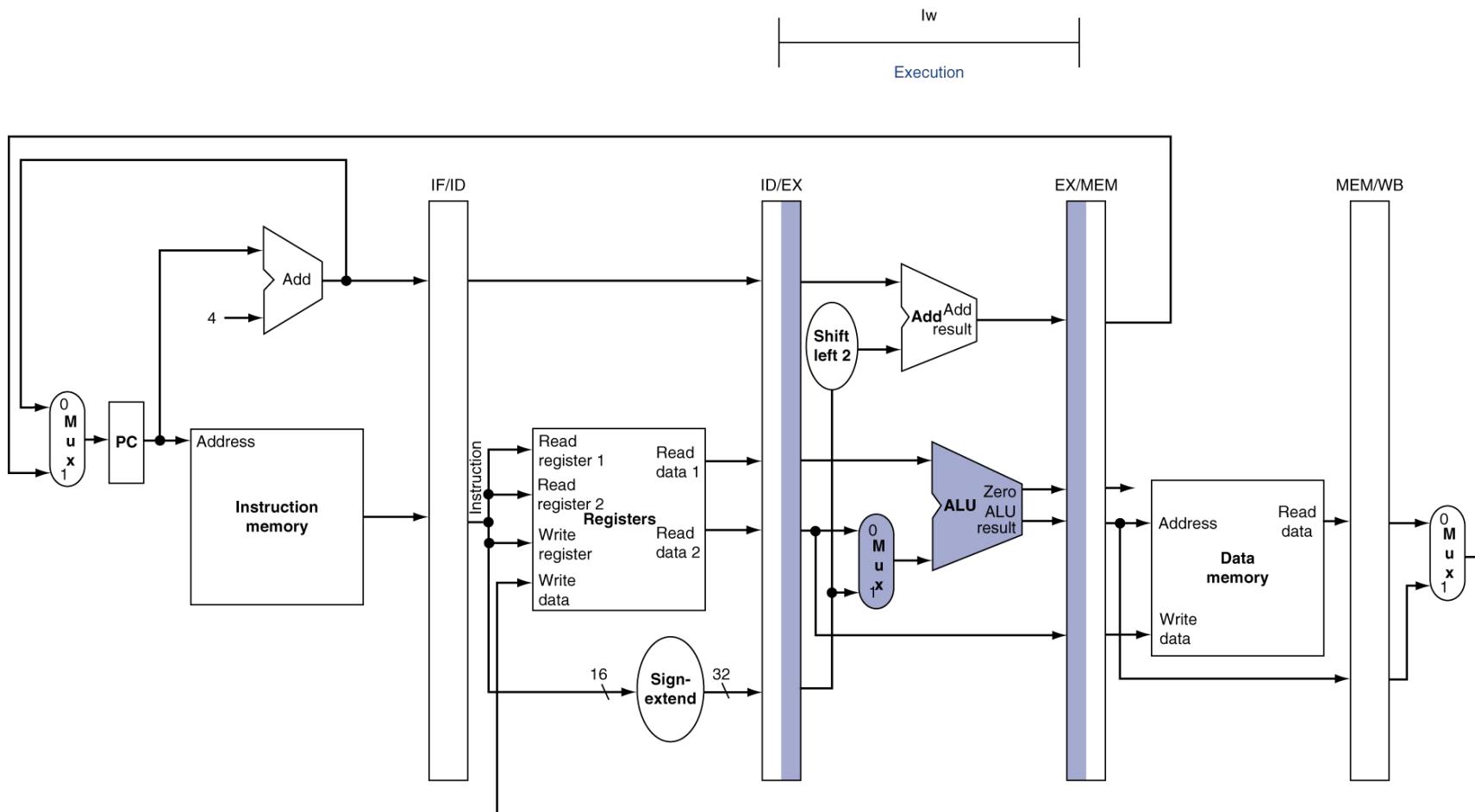
IF for Load, Store, ...



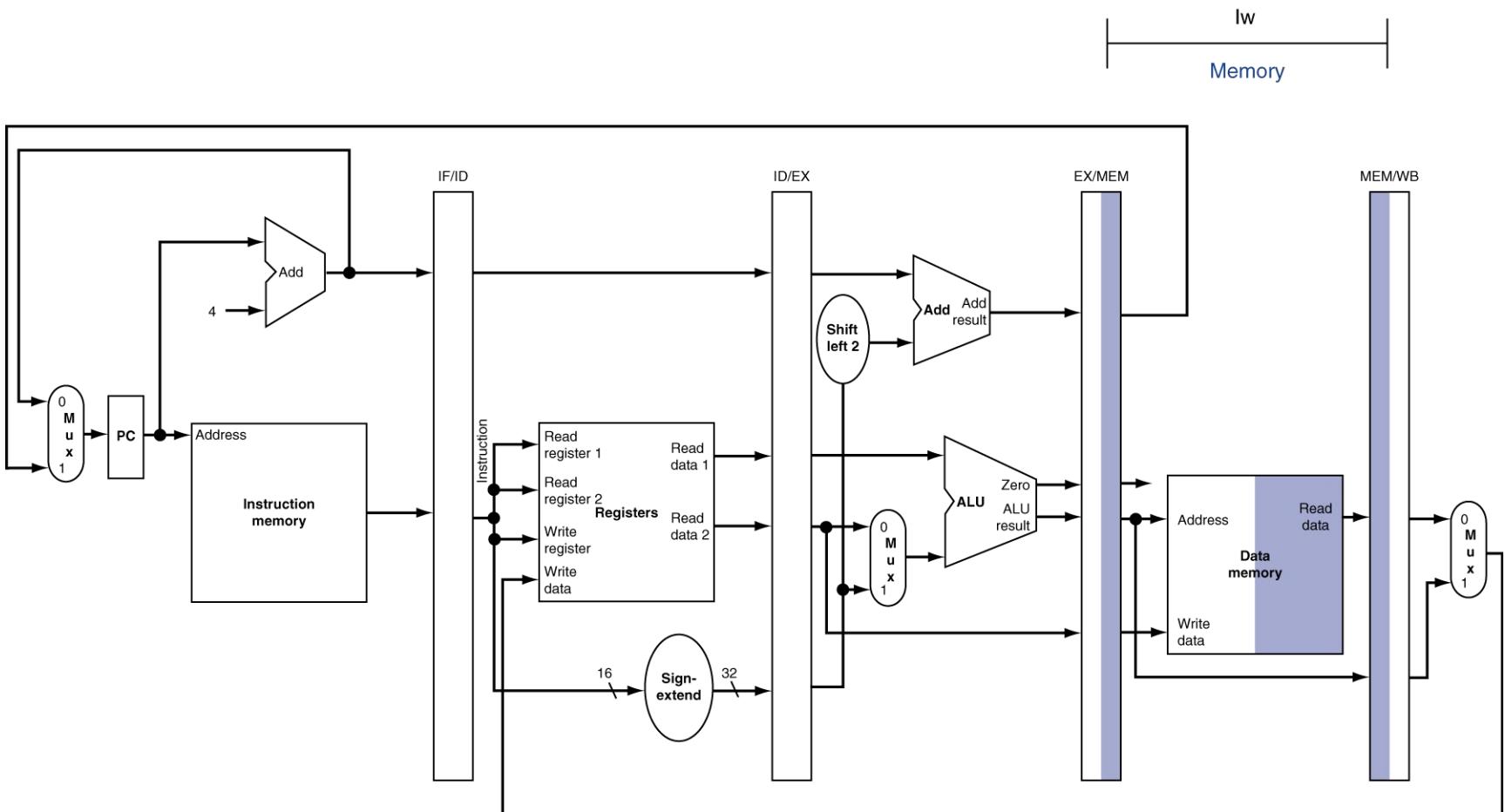
ID for Load, Store, ...



EX for Load

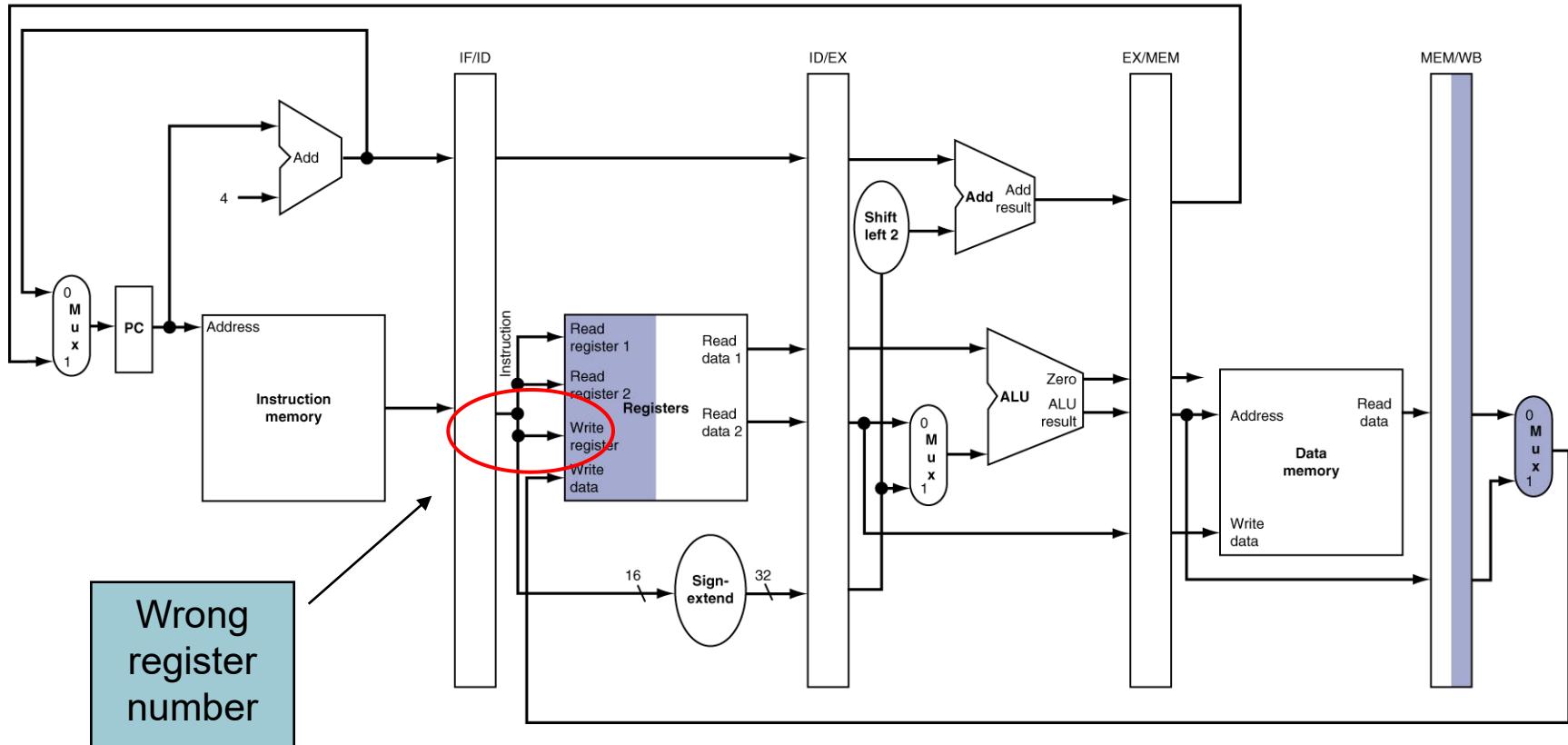


MEM for Load

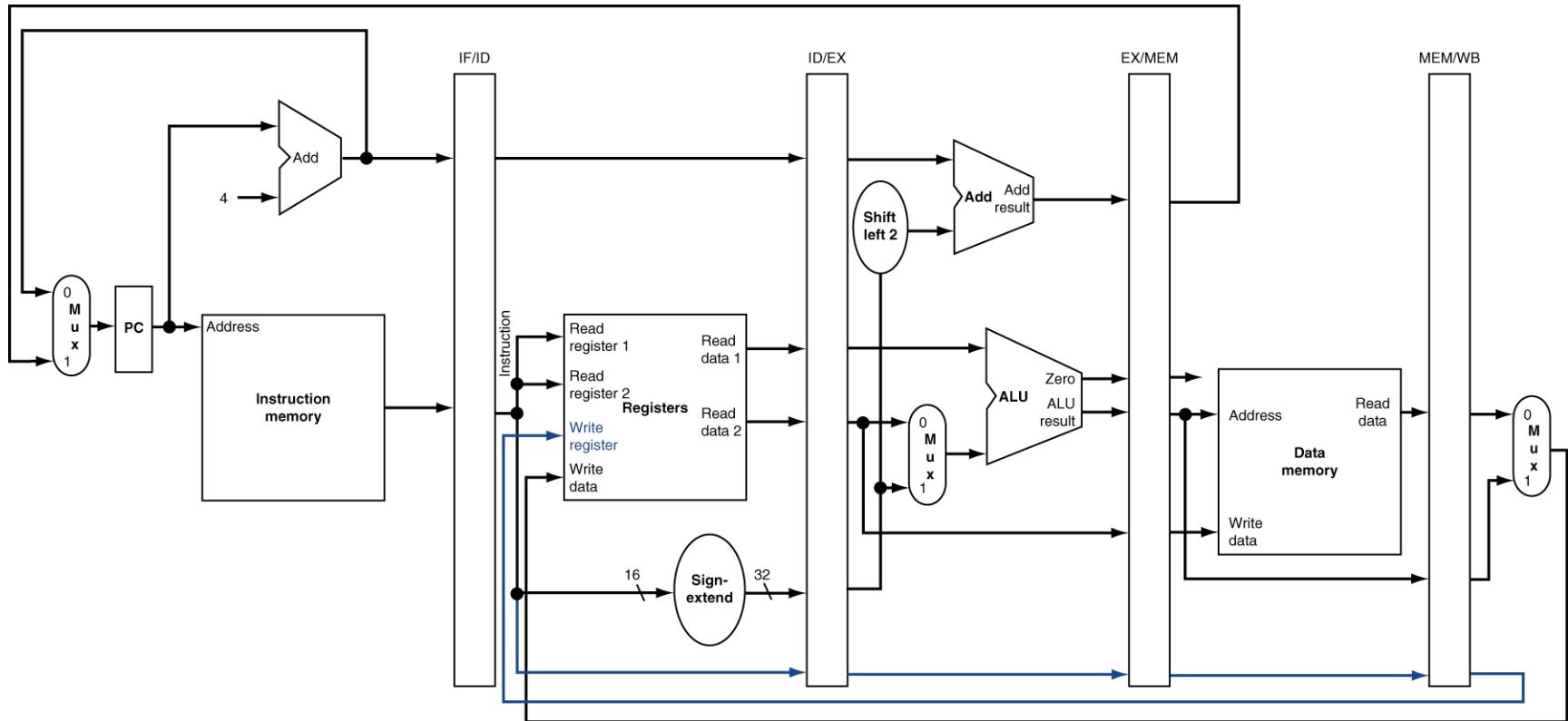


WB for Load

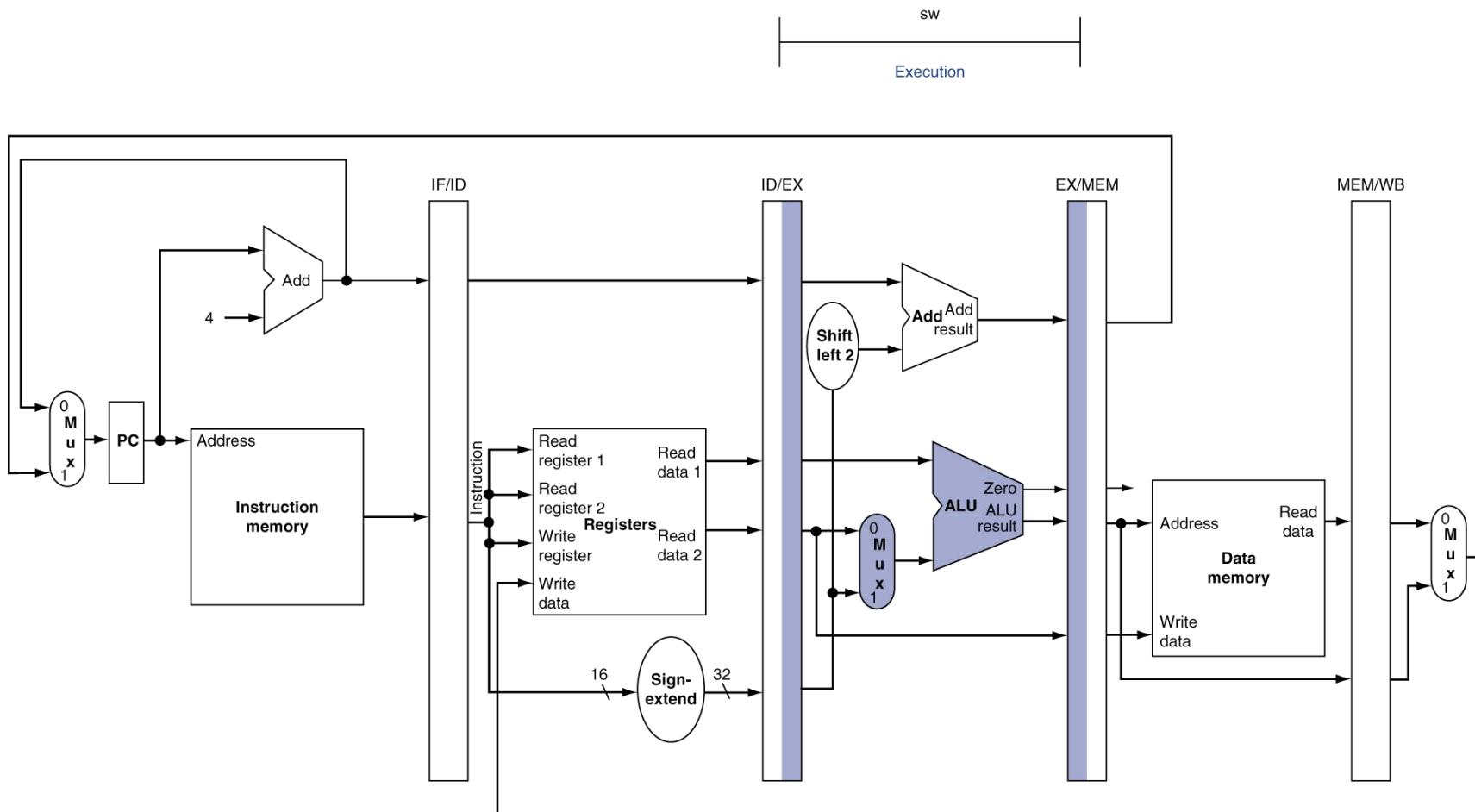
lw
Write back



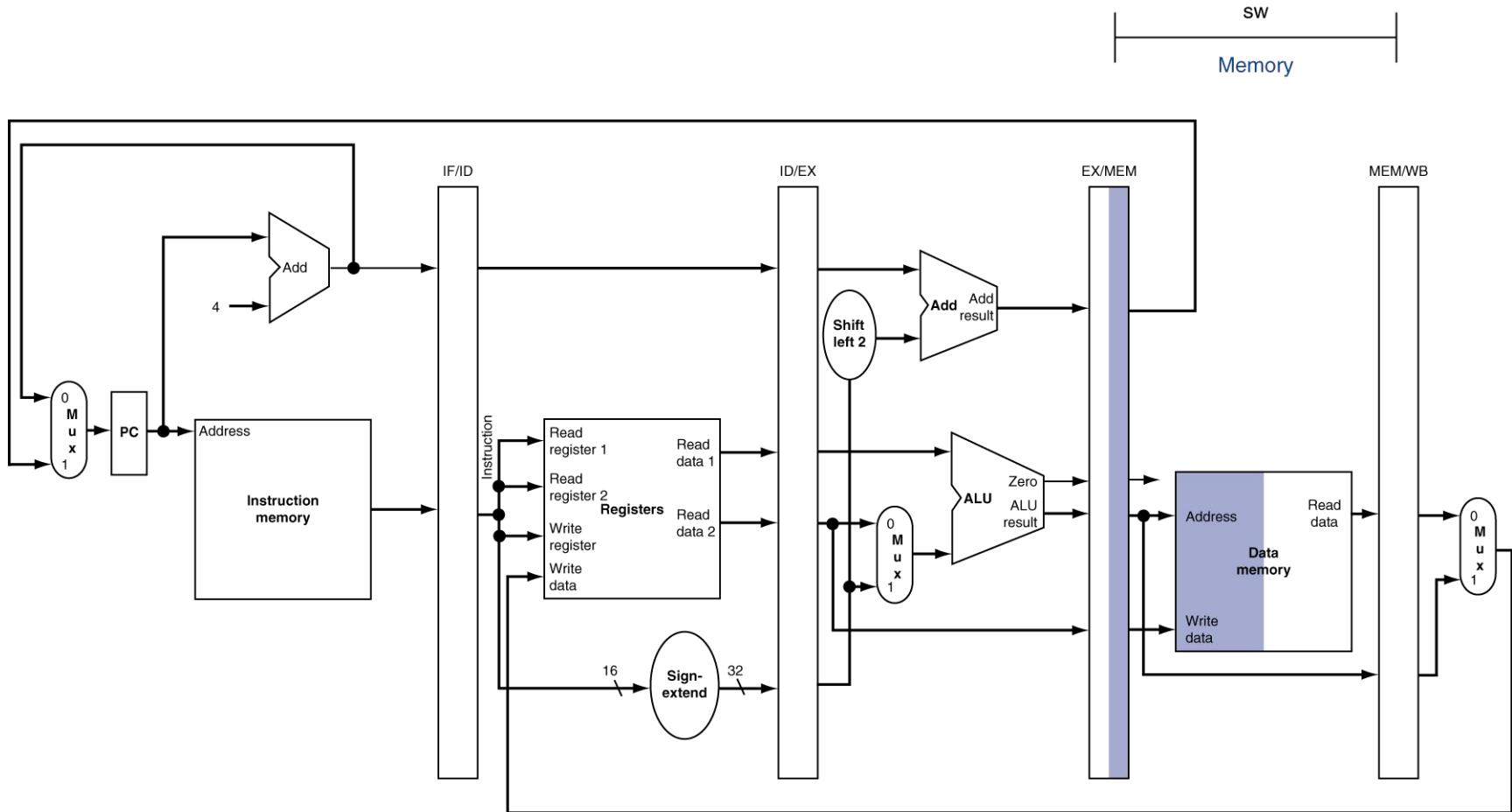
Corrected Datapath for Load



EX for Store

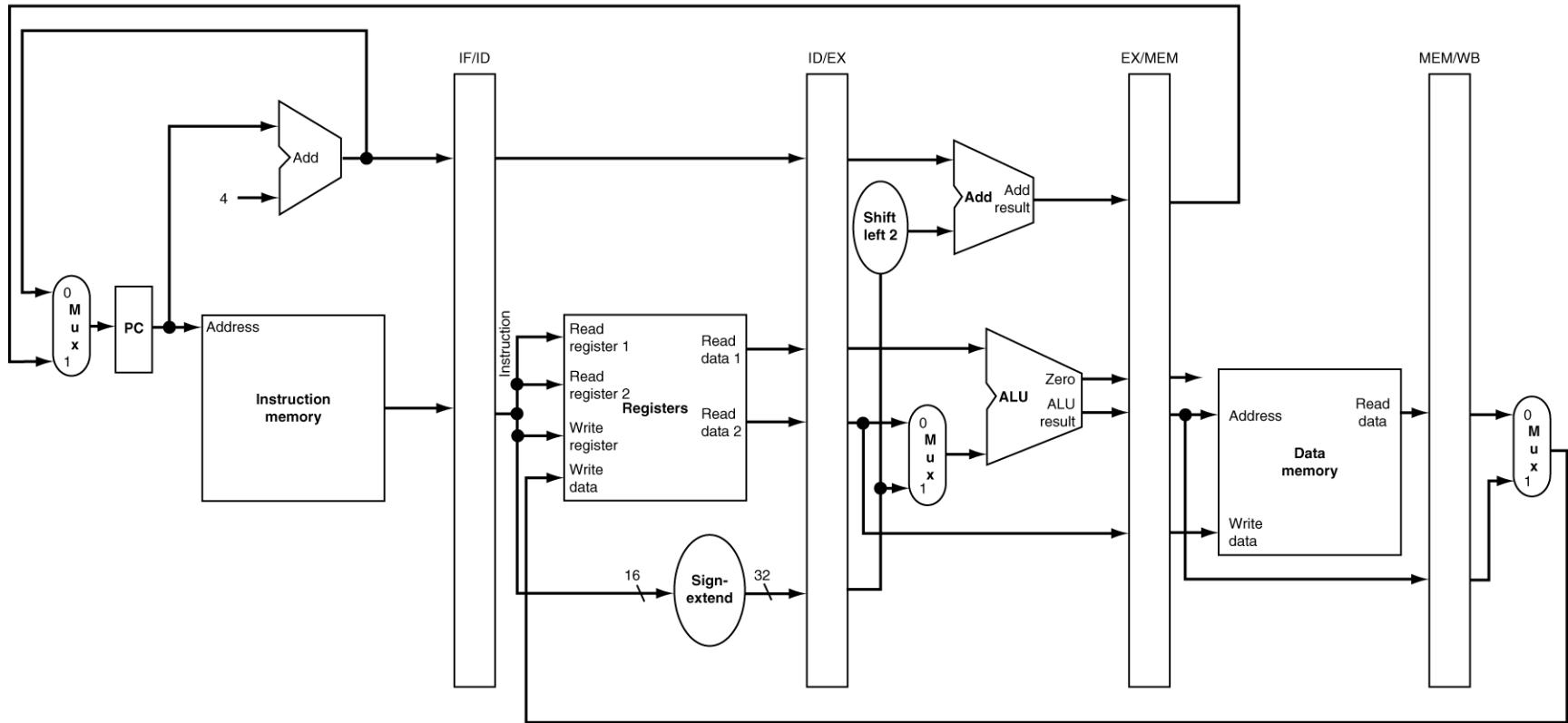


MEM for Store



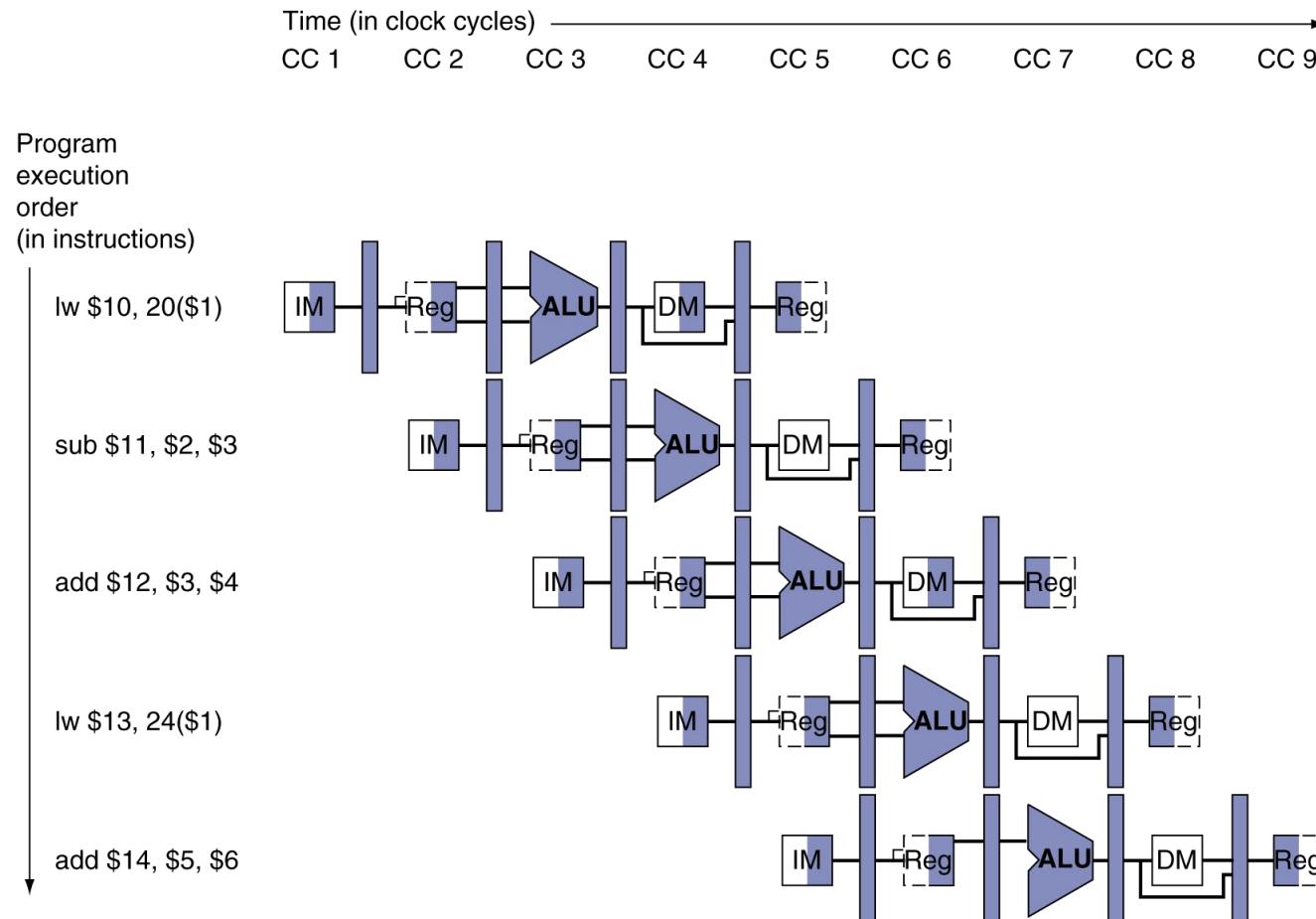
WB for Store

SW
Write-back



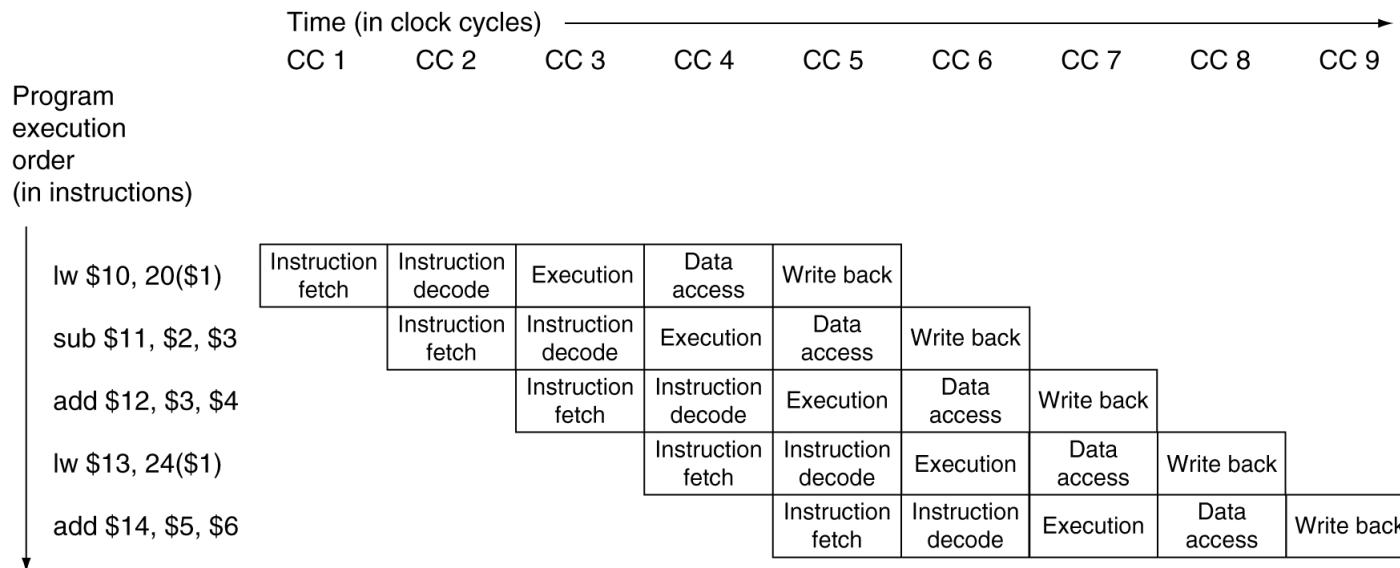
Multi-Cycle Pipeline Diagram

- Form showing resource usage



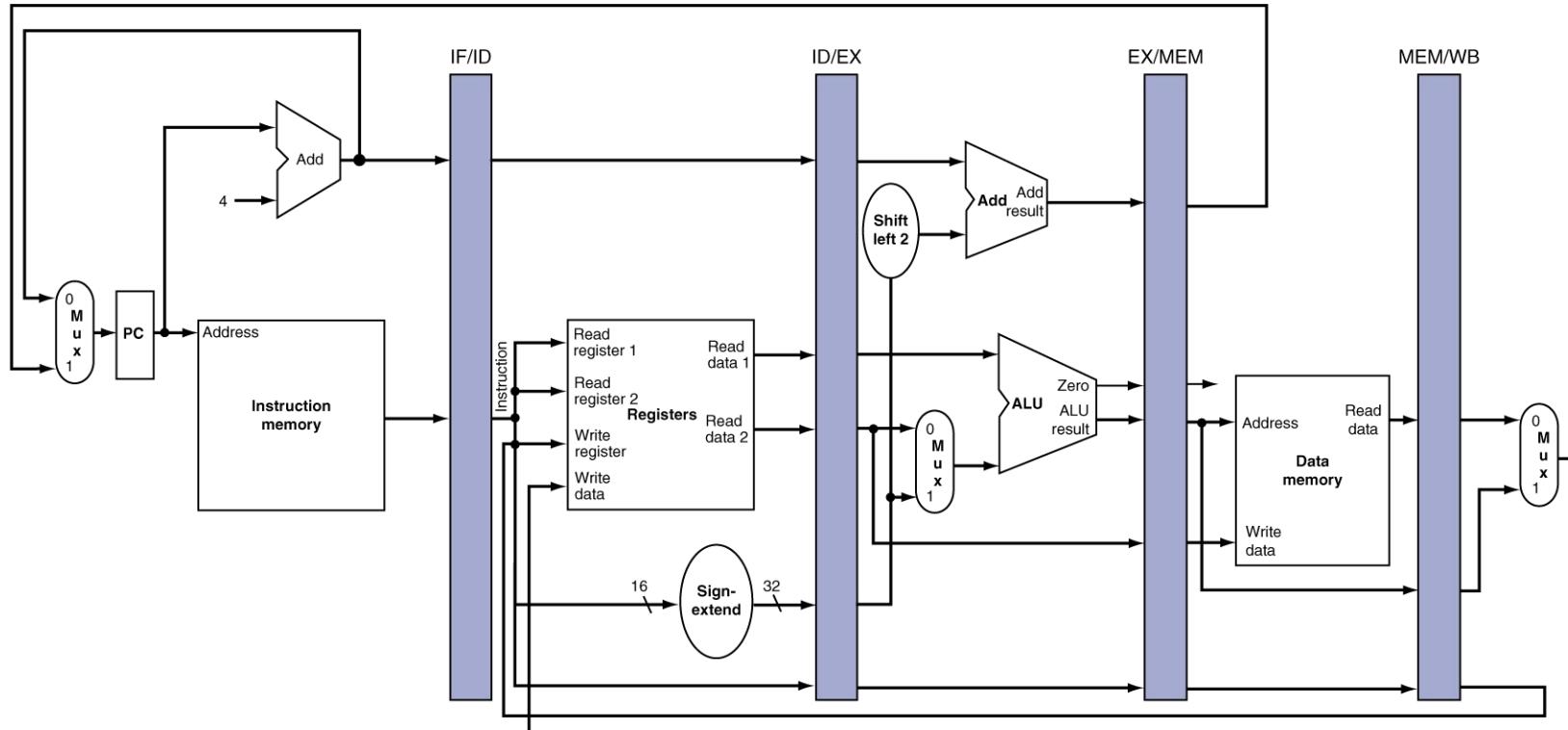
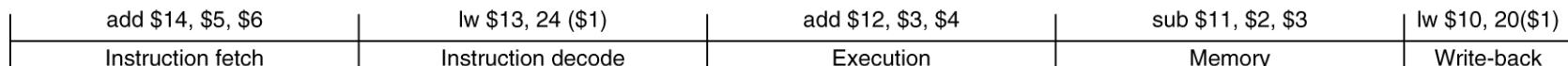
Multi-Cycle Pipeline Diagram

Traditional form

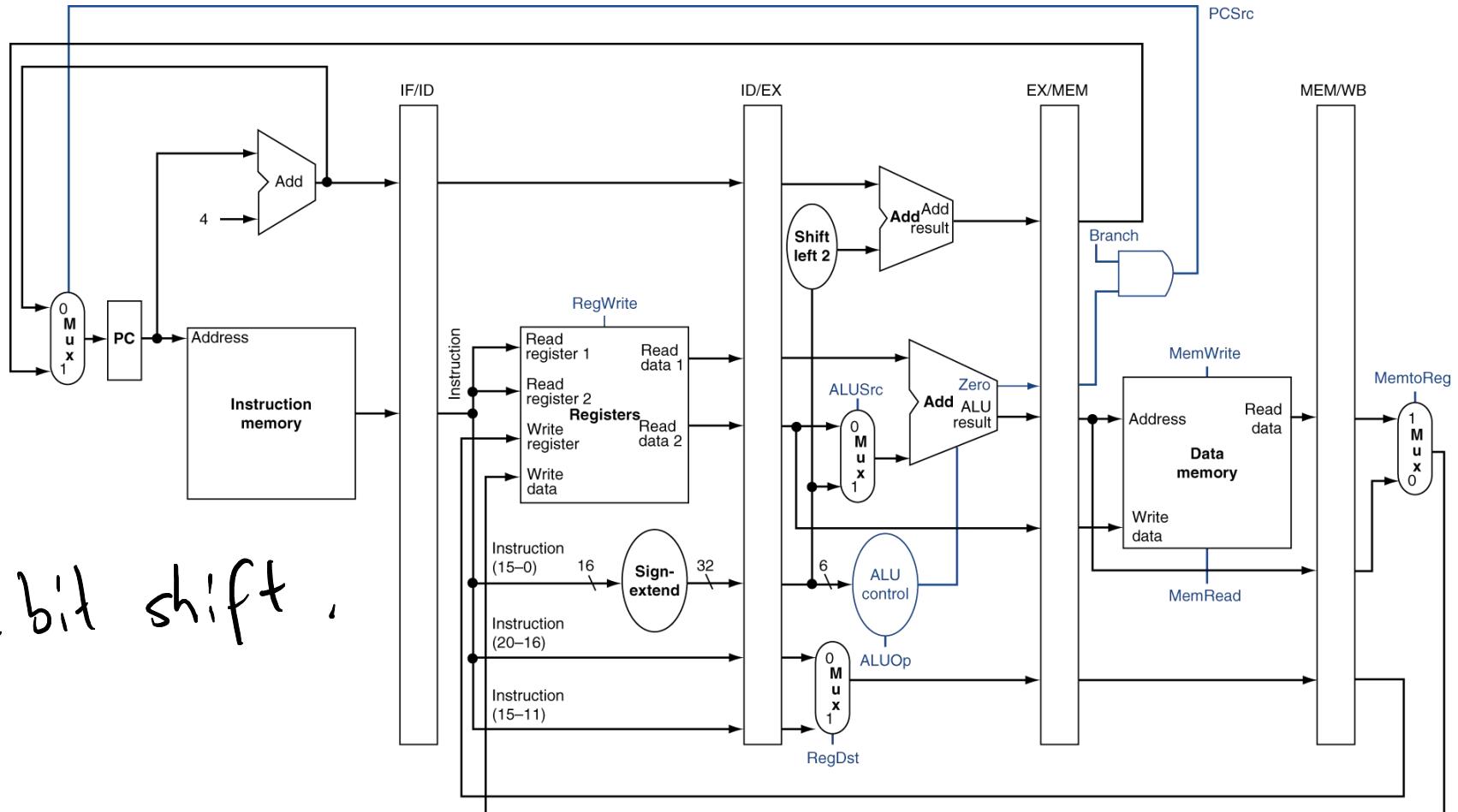


Single-Cycle Pipeline Diagram

State of pipeline in a given cycle

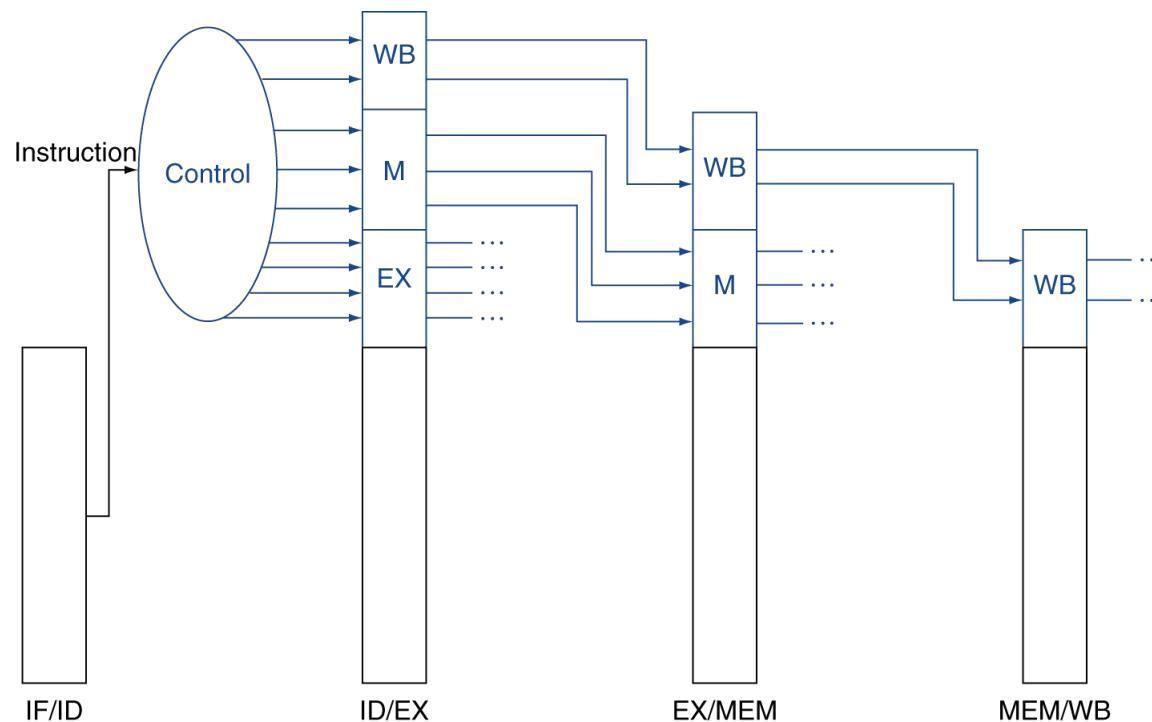


Pipelined Control (Simplified)

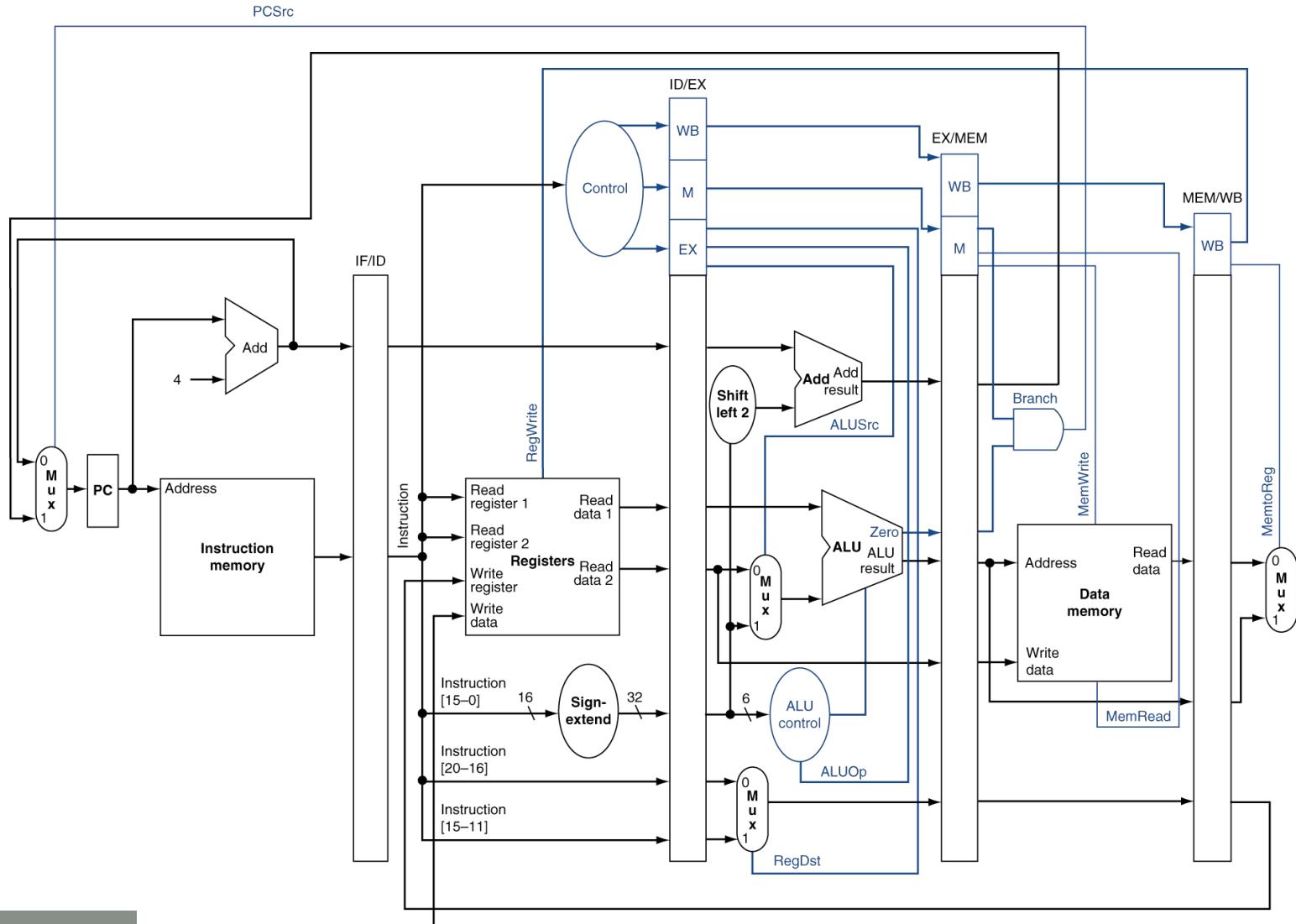


Pipelined Control

- Control signals derived from instruction
 - As in single-cycle implementation



Pipelined Control



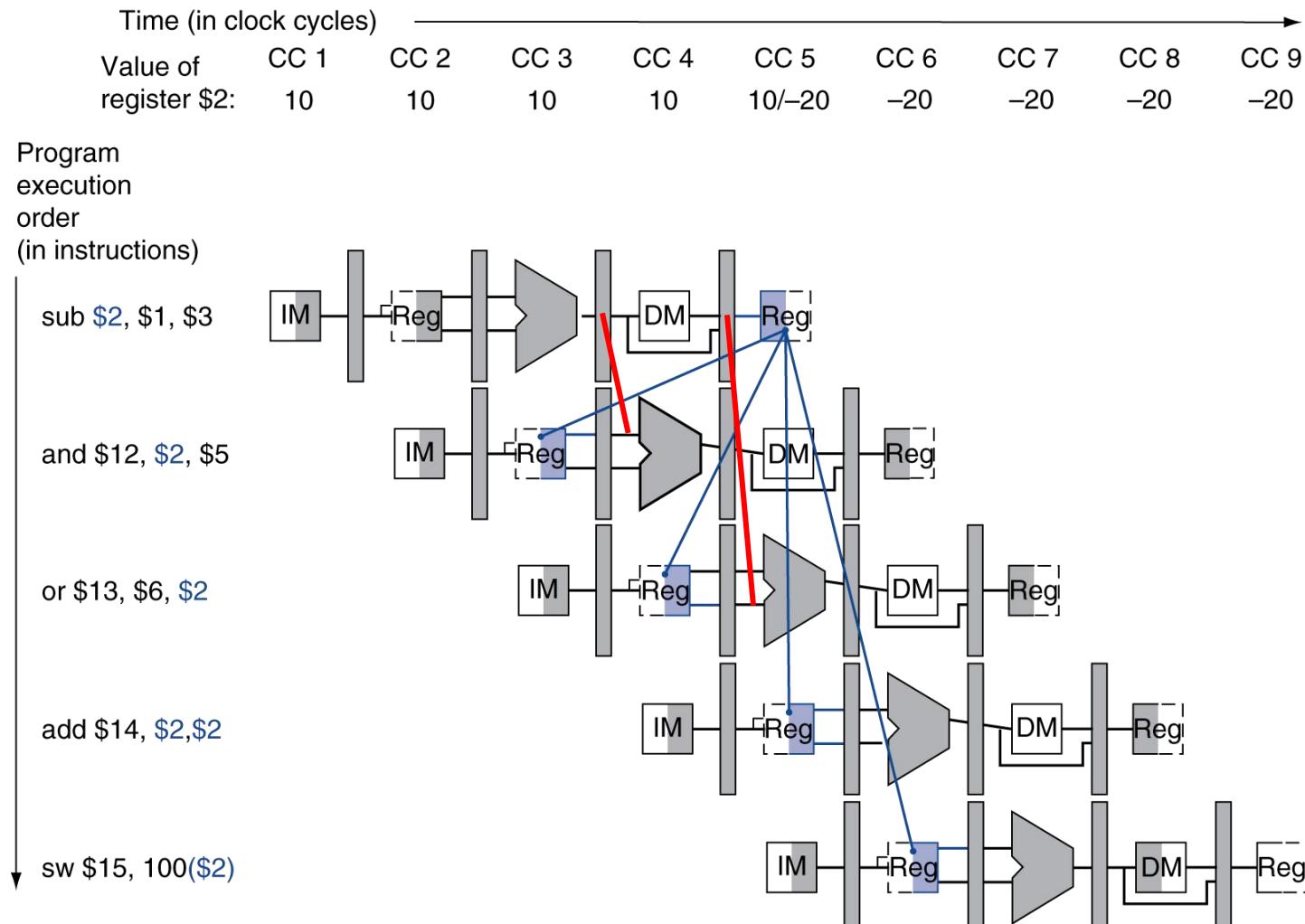
Data Hazards in ALU Instructions

- Consider this sequence:

```
sub $2, $1, $3  
and $12, $2, $5  
or $13, $6, $2  
add $14, $2, $2  
sw $15, 100($2)
```

- We can resolve hazards with forwarding
 - How do we detect when to forward?

Dependencies & Forwarding



Detecting the Need to Forward

- Pass register numbers along pipeline
 - e.g., ID/EX.RegisterRs = register number for Rs sitting in ID/EX pipeline register
- ALU operand register numbers in EX stage are given by
 - **ID/EX.RegisterRs, ID/EX.RegisterRt**
- Data hazards when
 - 1a. EX/MEM.RegisterRd = ID/EX.RegisterRs
 - 1b. EX/MEM.RegisterRd = ID/EX.RegisterRt
 - 2a. MEM/WB.RegisterRd = ID/EX.RegisterRs
 - 2b. MEM/WB.RegisterRd = ID/EX.RegisterRt

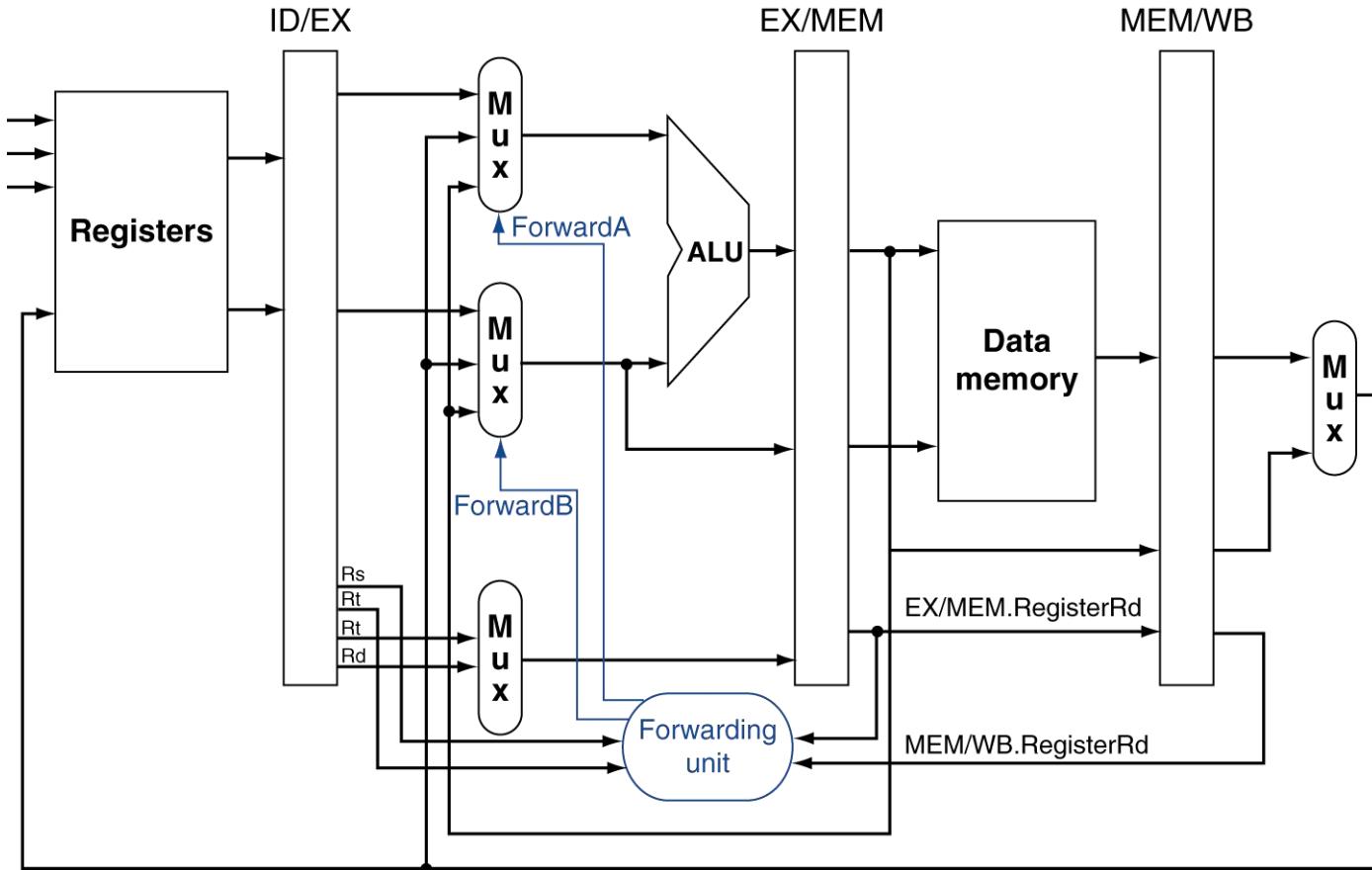
Fwd from
EX/MEM
pipeline reg

Fwd from
MEM/WB
pipeline reg

Detecting the Need to Forward

- But only if forwarding instruction will write to a register!
 - EX/MEM.RegWrite, MEM/WB.RegWrite
- And only if Rd for that instruction is not \$zero
 - EX/MEM.RegisterRd \neq 0,
MEM/WB.RegisterRd \neq 0

Forwarding Paths



b. With forwarding

Forwarding Conditions

- EX hazard
 - if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd = ID/EX.RegisterRs))
ForwardA = 10
 - if (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd = ID/EX.RegisterRt))
ForwardB = 10
- MEM hazard
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and (MEM/WB.RegisterRd = ID/EX.RegisterRs))
ForwardA = 01
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and (MEM/WB.RegisterRd = ID/EX.RegisterRt))
ForwardB = 01

Double Data Hazard

- Consider the sequence:

add \$1, \$1, \$2

add \$1, \$1, \$3

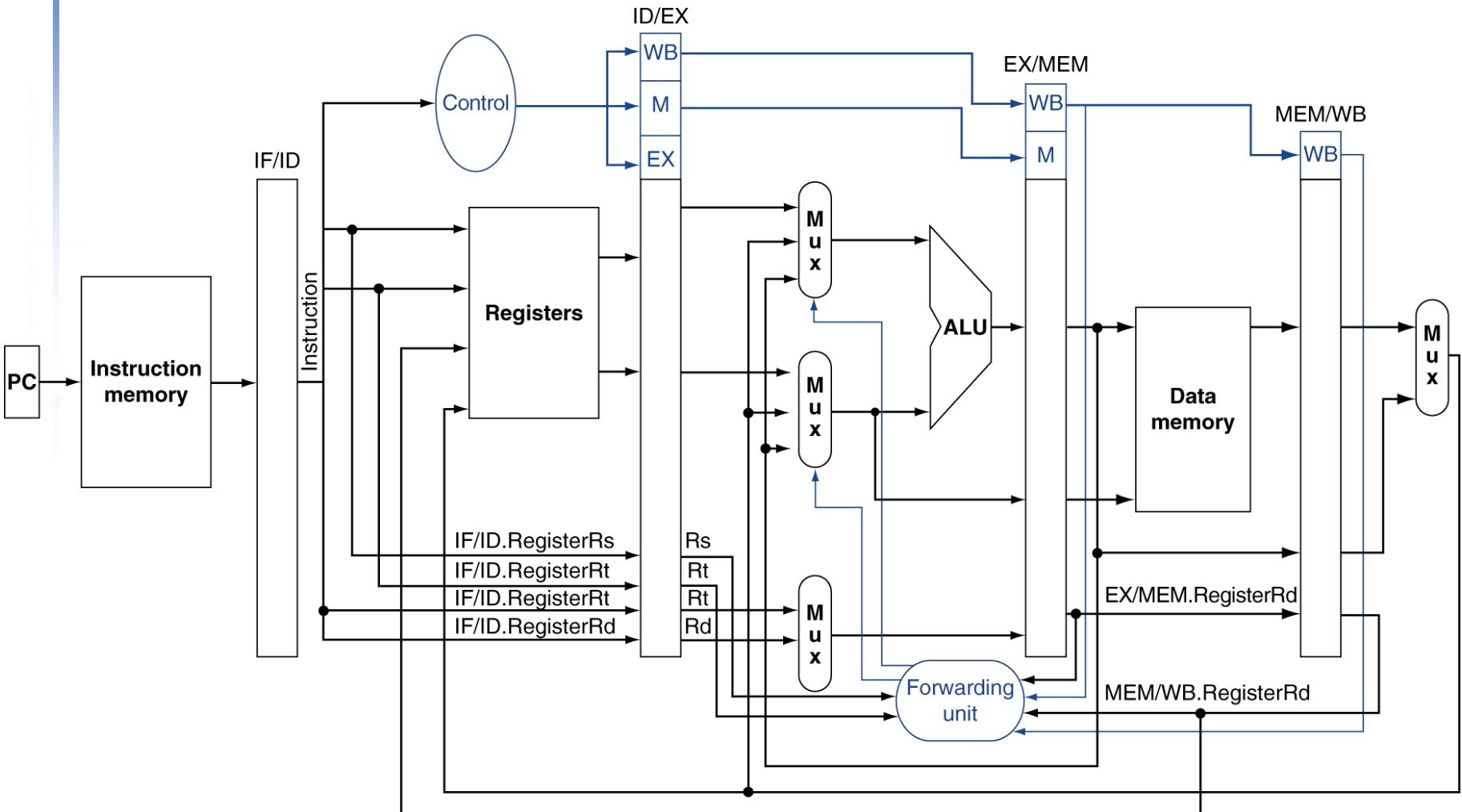
add \$1, \$1, \$4

- Both hazards occur
 - Want to use the most recent
- Revise MEM hazard condition
 - Only fwd if EX hazard condition isn't true

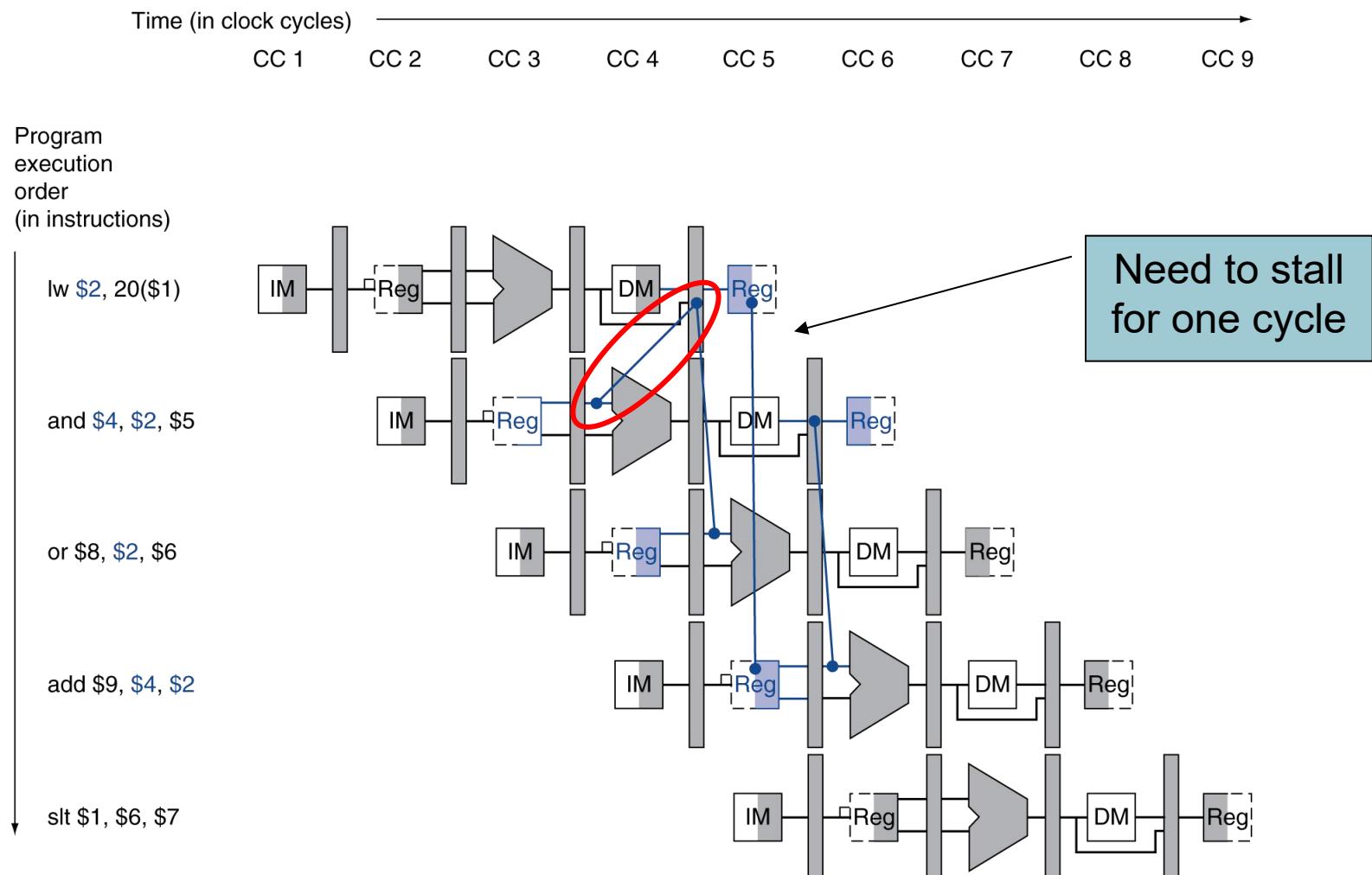
Revised Forwarding Condition

- MEM hazard
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd = ID/EX.RegisterRs))
and (MEM/WB.RegisterRd = ID/EX.RegisterRs))
ForwardA = 01
 - if (MEM/WB.RegWrite and (MEM/WB.RegisterRd ≠ 0)
and not (EX/MEM.RegWrite and (EX/MEM.RegisterRd ≠ 0)
and (EX/MEM.RegisterRd = ID/EX.RegisterRt))
and (MEM/WB.RegisterRd = ID/EX.RegisterRt))
ForwardB = 01

Datapath with Forwarding



Load-Use Data Hazard



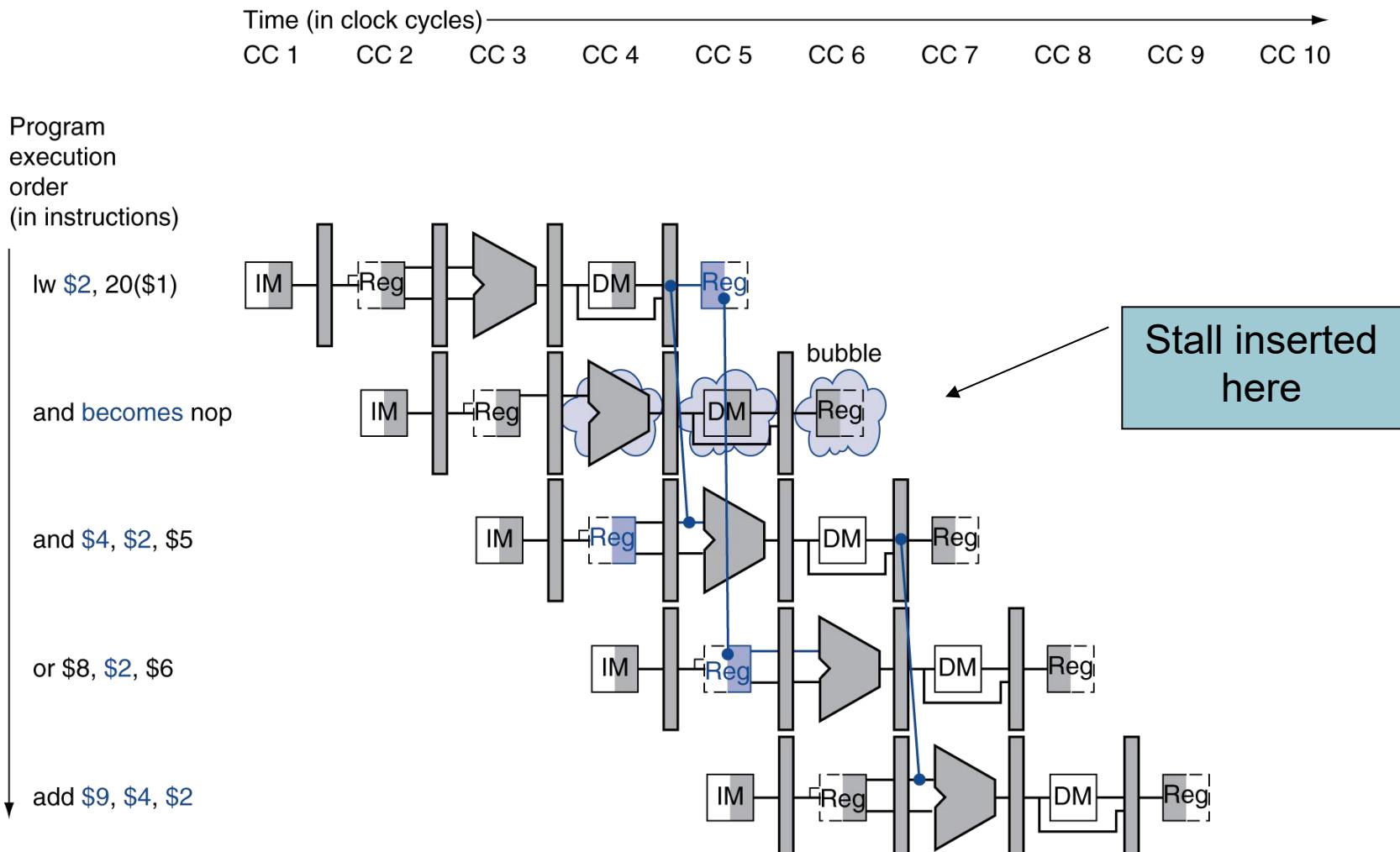
Load-Use Hazard Detection

- Check when using instruction is decoded in ID stage
- ALU operand register numbers in ID stage are given by
 - IF/ID.RegisterRs, IF/ID.RegisterRt
- Load-use hazard when
 - ID/EX.MemRead and
$$((ID/EX.RegisterRt = IF/ID.RegisterRs) \text{ or } (ID/EX.RegisterRt = IF/ID.RegisterRt))$$
- If detected, stall and insert bubble

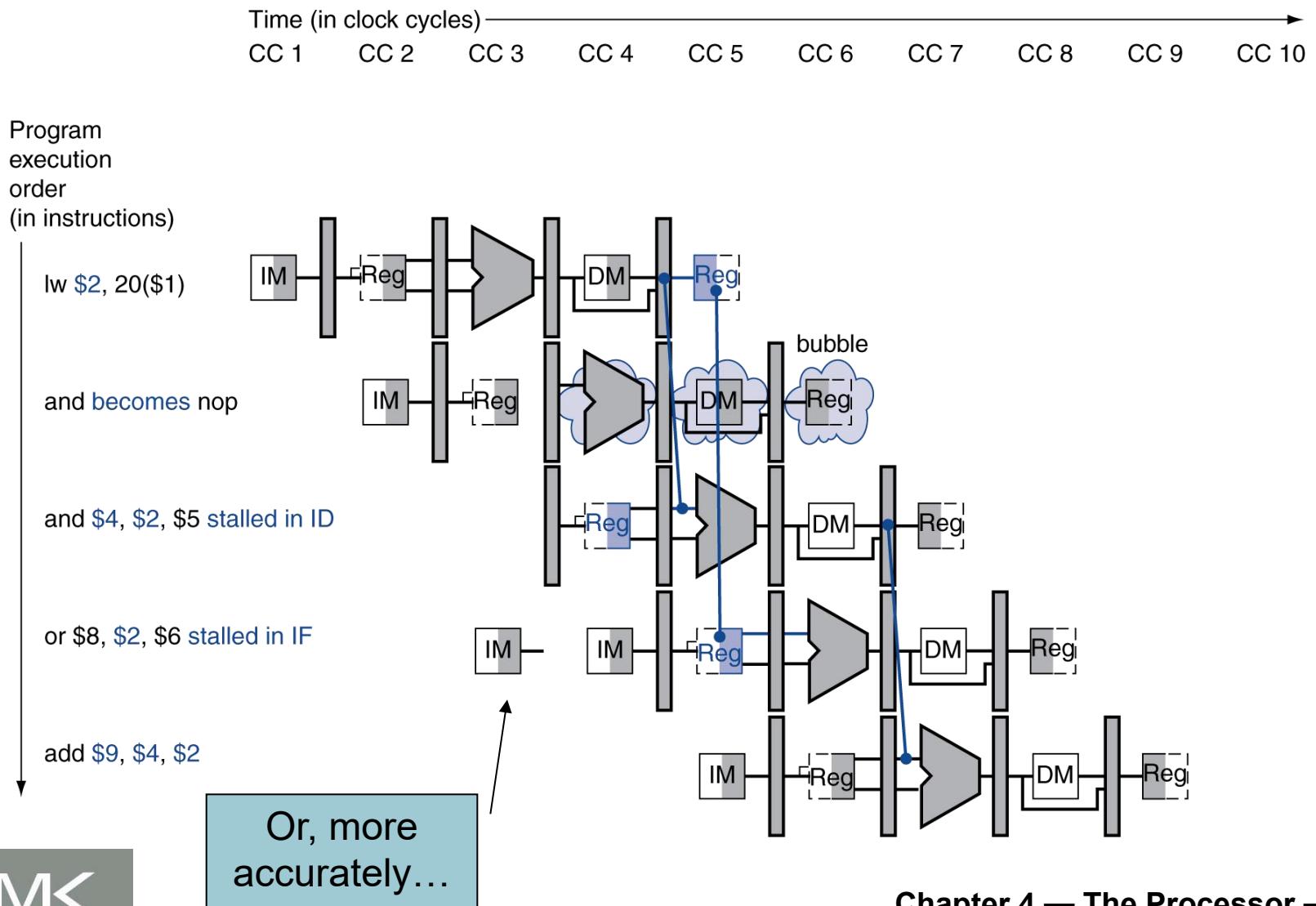
How to Stall the Pipeline

- Force control values in ID/EX register to 0
 - EX, MEM and WB do nop (no-operation)
- Prevent update of PC and IF/ID register
 - Using instruction is decoded again
 - Following instruction is fetched again
 - 1-cycle stall allows MEM to read data for 1w
 - Can subsequently forward to EX stage

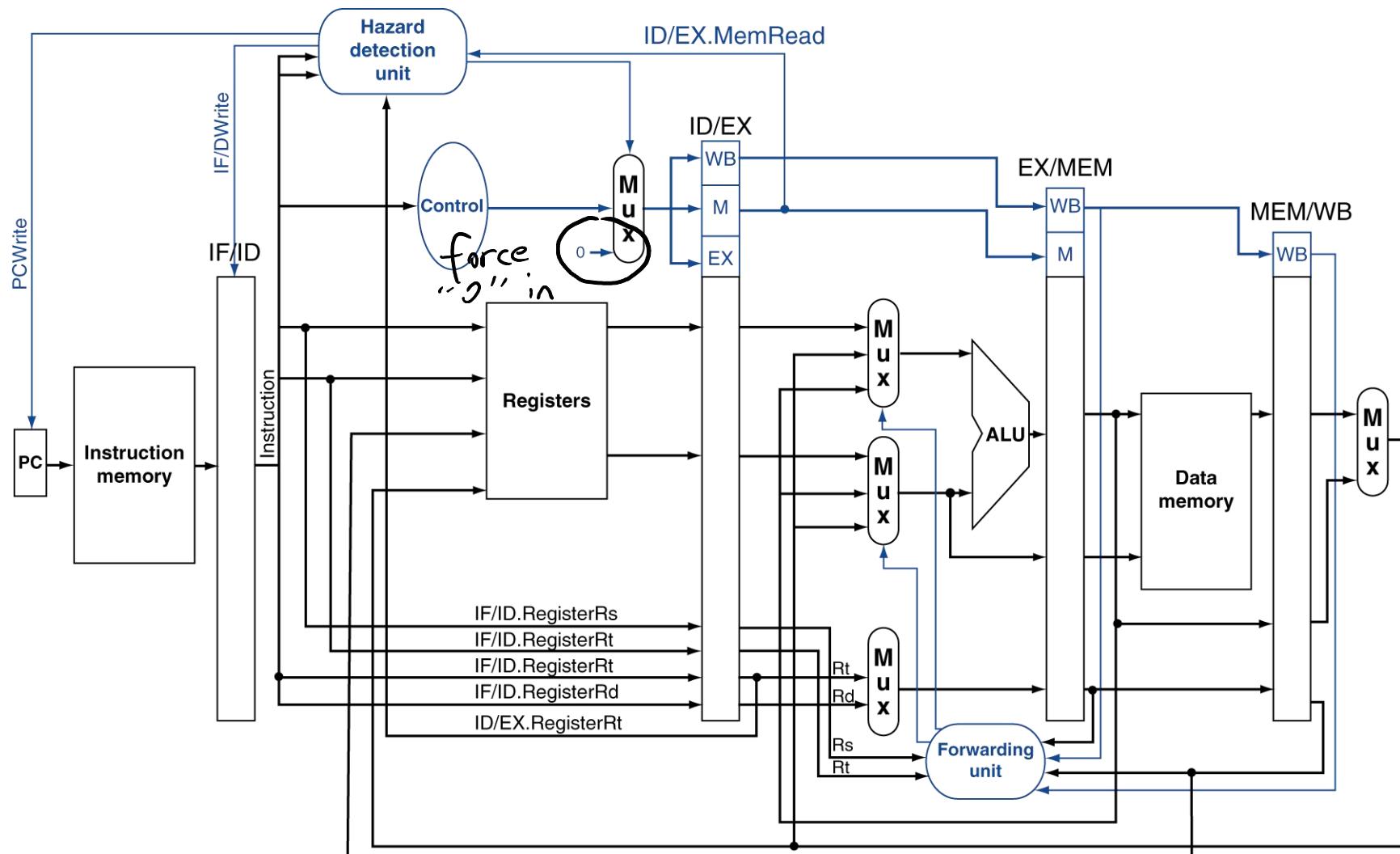
Stall/Bubble in the Pipeline



Stall/Bubble in the Pipeline



Datapath with Hazard Detection



Stalls and Performance

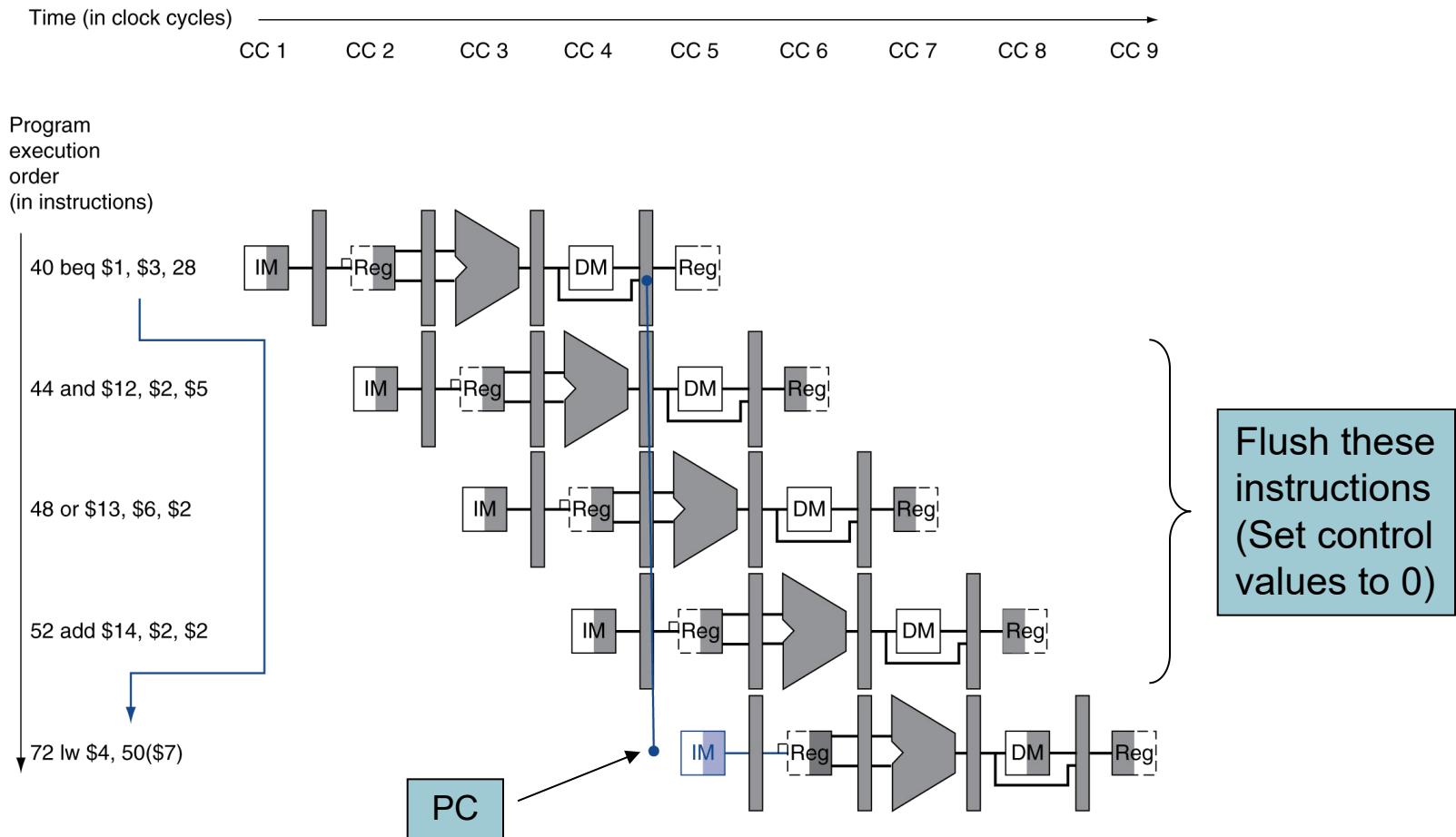
The BIG Picture

Register 0 is always 0.

- Stalls reduce performance
 - But are required to get correct results
- Compiler can arrange code to avoid hazards and stalls
 - Requires knowledge of the pipeline structure

Branch Hazards

- If branch outcome determined in MEM

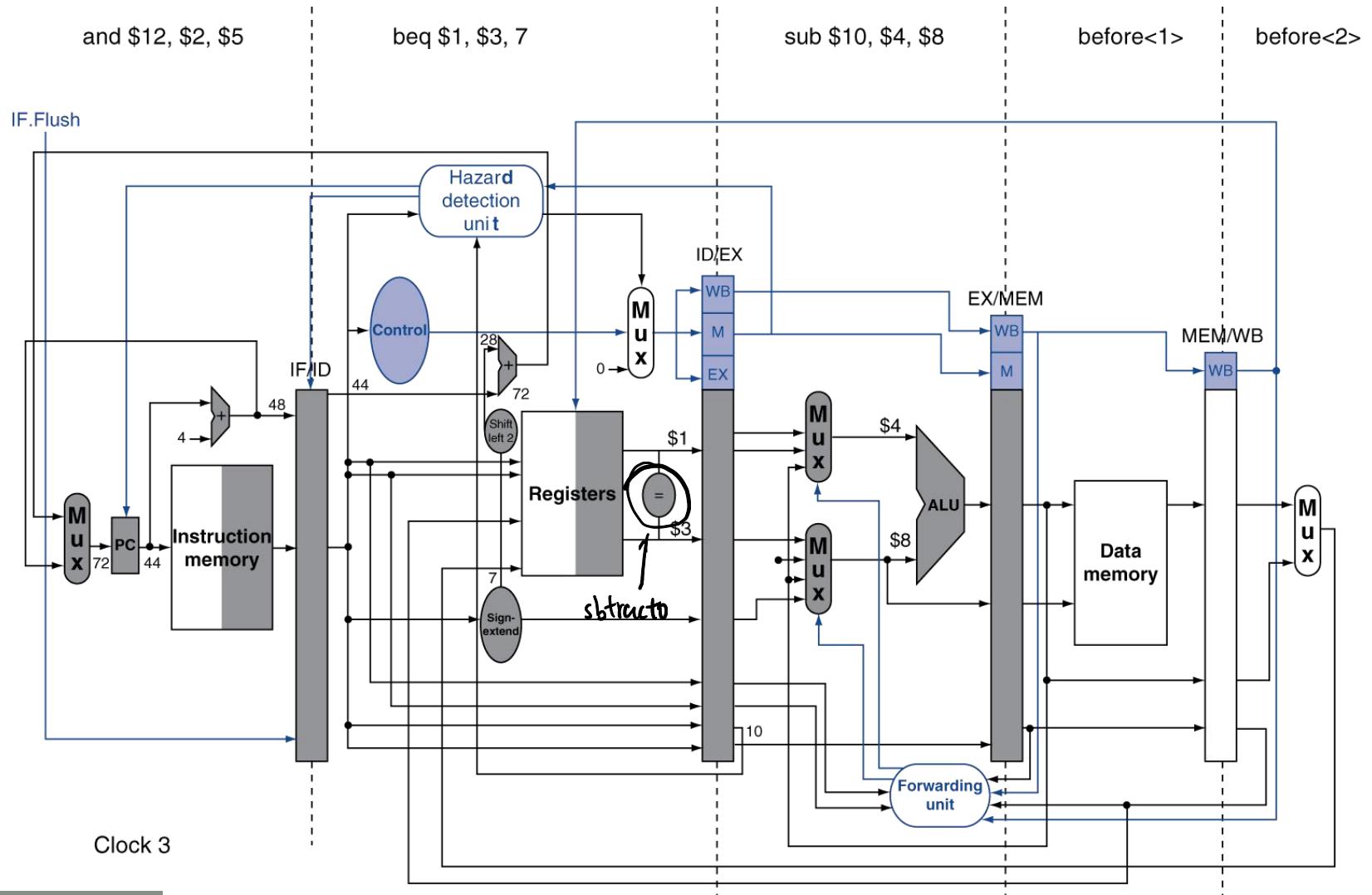


Reducing Branch Delay

- Move hardware to determine outcome to ID stage
 - Target address adder
 - Register comparator
- Example: branch taken

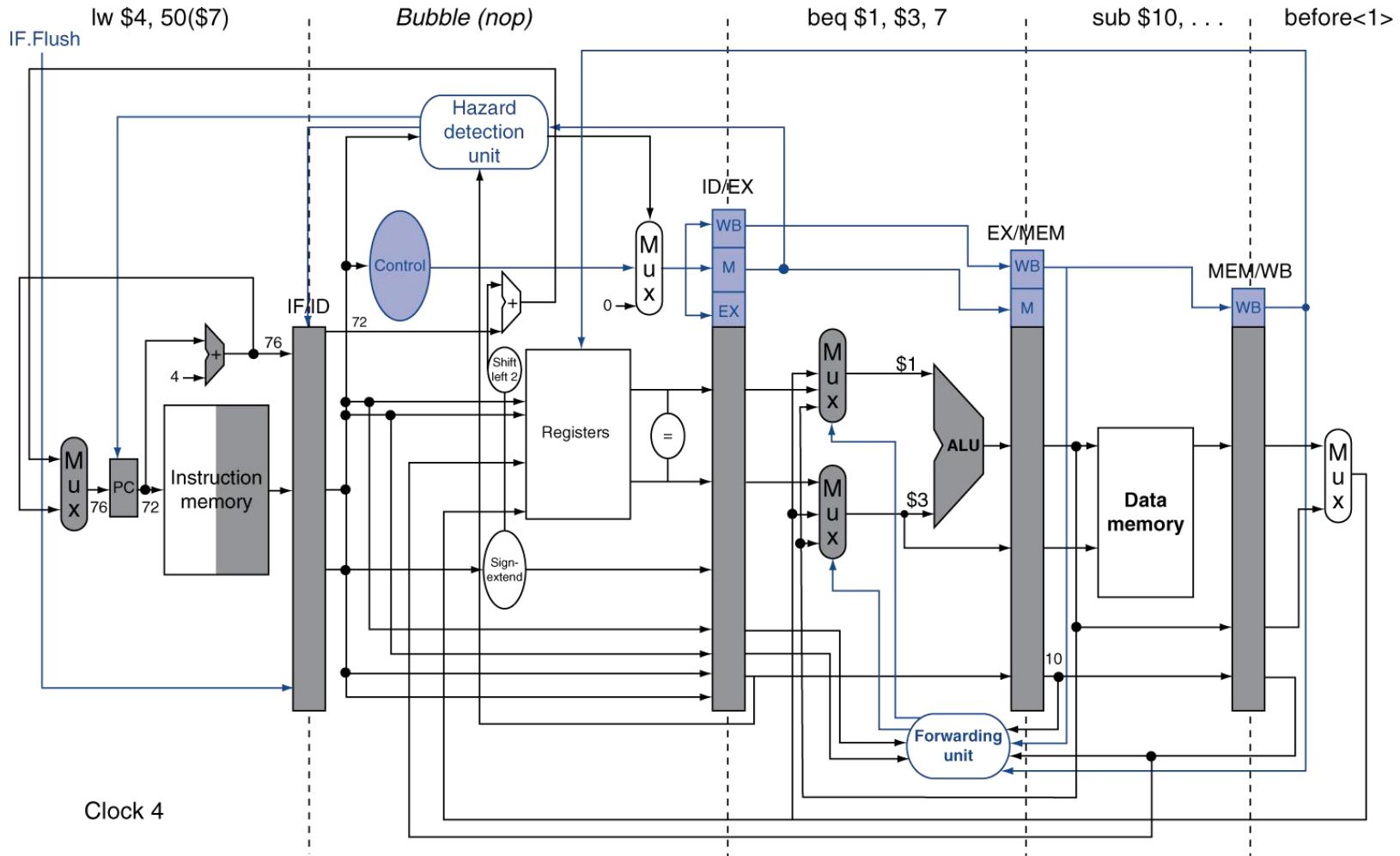
```
36:    sub   $10, $4, $8  
40:    beq   $1,  $3,  7  
44:    and   $12, $2, $5  
48:    or    $13, $2, $6  
52:    add   $14, $4, $2  
56:    slt   $15, $6, $7  
      ...  
72:    lw    $4, 50($7)
```

Example: Branch Taken



Clock 3

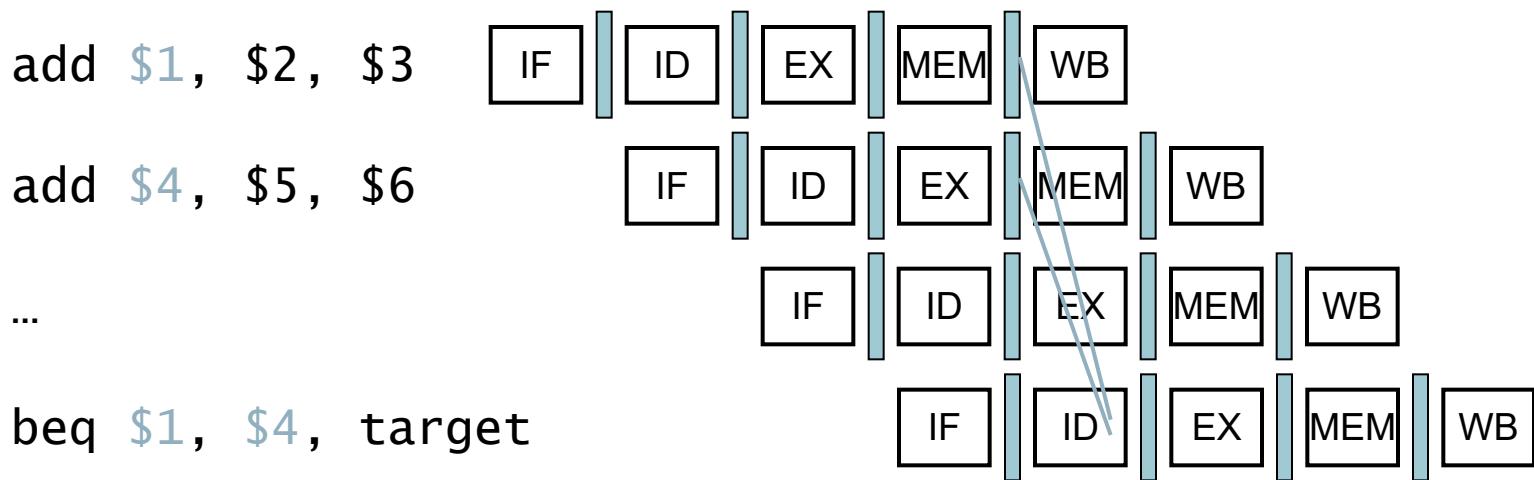
Example: Branch Taken



Clock 4

Data Hazards for Branches

- If a comparison register is a destination of 2nd or 3rd preceding ALU instruction

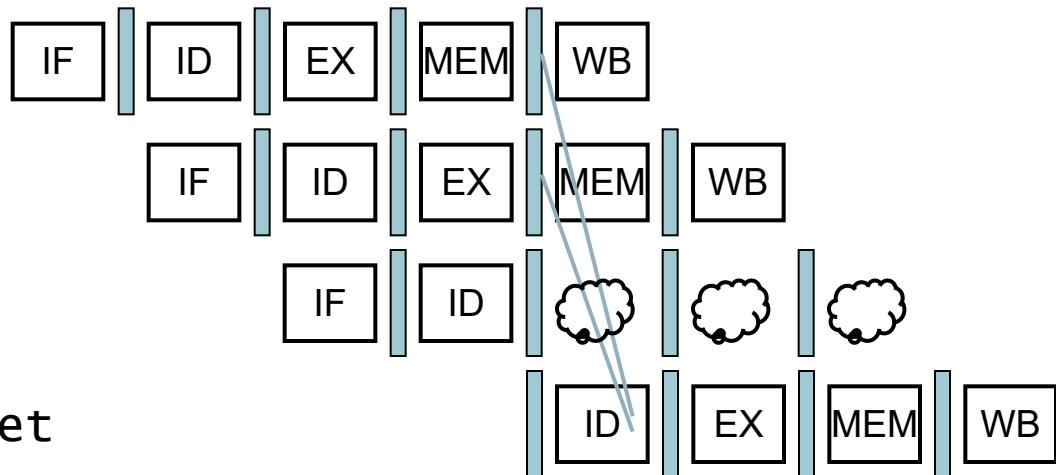


- Can resolve using forwarding

Data Hazards for Branches

- If a comparison register is a destination of preceding ALU instruction or 2nd preceding load instruction
 - Need 1 stall cycle

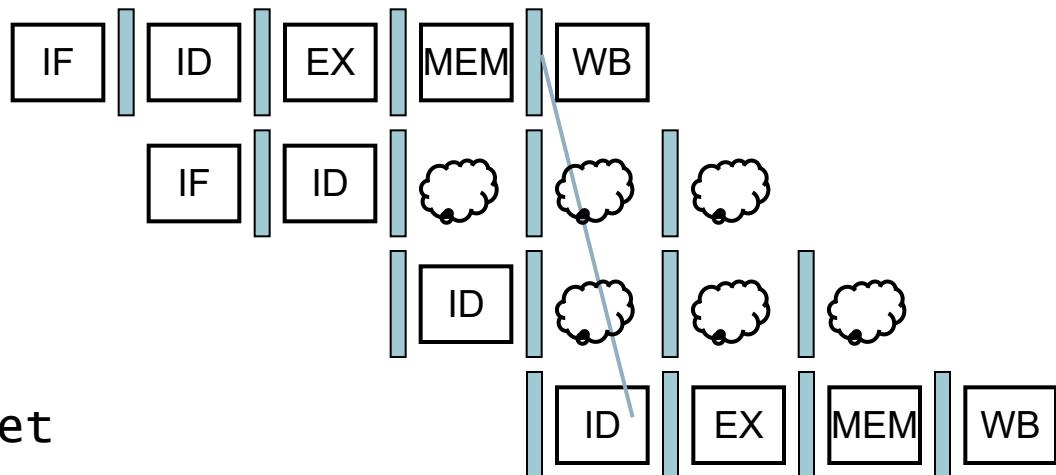
lw \$1, addr
add \$4, \$5, \$6
beq stalled
beq \$1, \$4, target



Data Hazards for Branches

- If a comparison register is a destination of immediately preceding load instruction
 - Need 2 stall cycles

lw \$1, addr
beq stalled
beq stalled
beq \$1, \$0, target

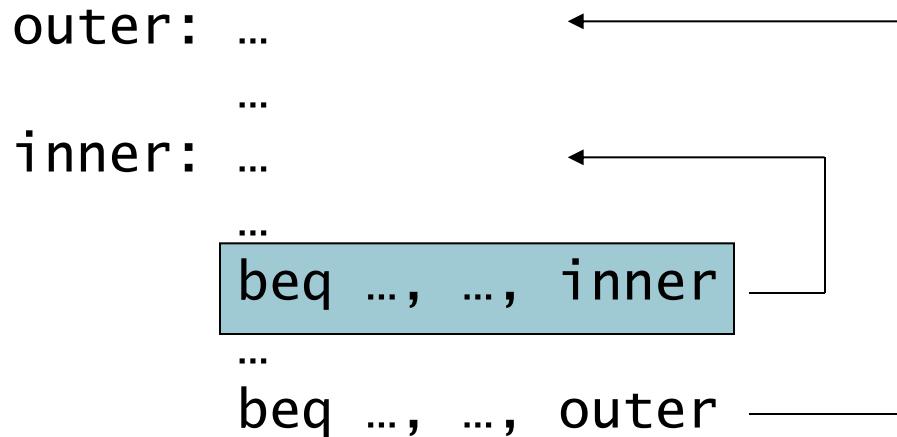


Dynamic Branch Prediction

- In deeper and superscalar pipelines, branch penalty is more significant
- Use dynamic prediction
 - Branch prediction buffer (aka branch history table)
 - Indexed by recent branch instruction addresses
 - Stores outcome (taken/not taken)
 - To execute a branch
 - Check table, expect the same outcome
 - Start fetching from fall-through or target
 - If wrong, flush pipeline and flip prediction

1-Bit Predictor: Shortcoming

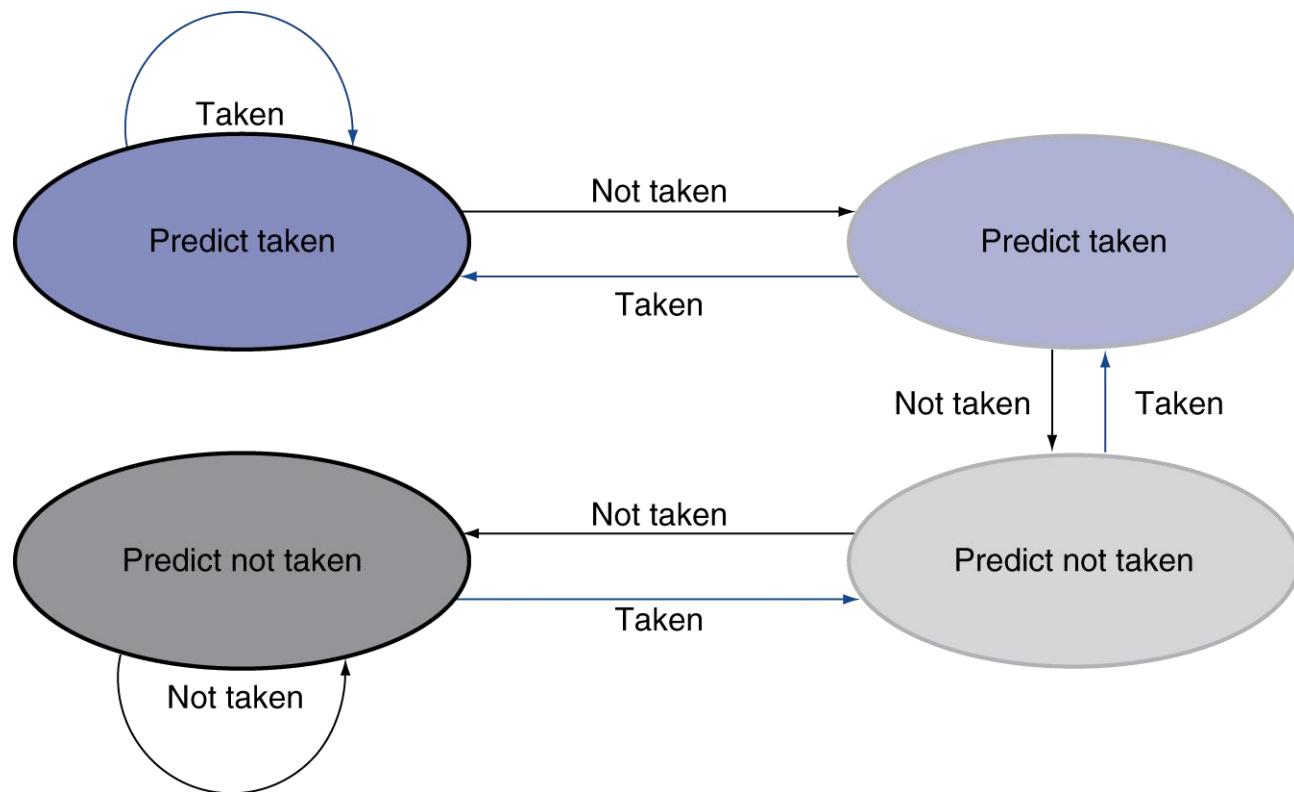
- Inner loop branches mispredicted twice!



- Mispredict as taken on last iteration of inner loop
- Then mispredict as not taken on first iteration of inner loop next time around

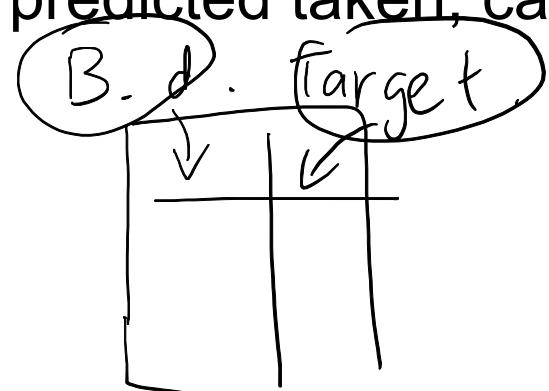
2-Bit Predictor

- Only change prediction on two successive mispredictions



Calculating the Branch Target

- Even with predictor, still need to calculate the target address
 - 1-cycle penalty for a taken branch
- **Branch target buffer** → Branch address
 - Cache of target addresses and Target .
 - Indexed by PC when instruction fetched
 - If hit and instruction is branch predicted taken, can fetch target immediately



Exceptions and Interrupts

e.g., Arithmetic

- “Unexpected” events requiring change in flow of control
 - Target Address outside Program allocated mem
 - Different ISAs use the terms differently
- Exception
 - Arises within the CPU
 - e.g., undefined opcode, overflow, syscall, ...
- Interrupt
 - Bad update
 - From an external I/O controller
- Dealing with them without sacrificing performance is hard

Handling Exceptions

- In MIPS, exceptions managed by a System Control Coprocessor (CP0)
- Save PC of offending (or interrupted) instruction
 - In MIPS: Exception Program Counter (EPC)
- Save indication of the problem
 - In MIPS: Cause register
 - We'll assume 1-bit
 - 0 for undefined opcode, 1 for overflow
- Jump to handler at 8000 00180

An Alternate Mechanism

- Vectored Interrupts
 - Handler address determined by the cause
- Example:
 - Undefined opcode: C000 0000
 - Overflow: C000 0020
 - C000 0040
- Instructions either
 - Deal with the interrupt, or
 - Jump to real handler

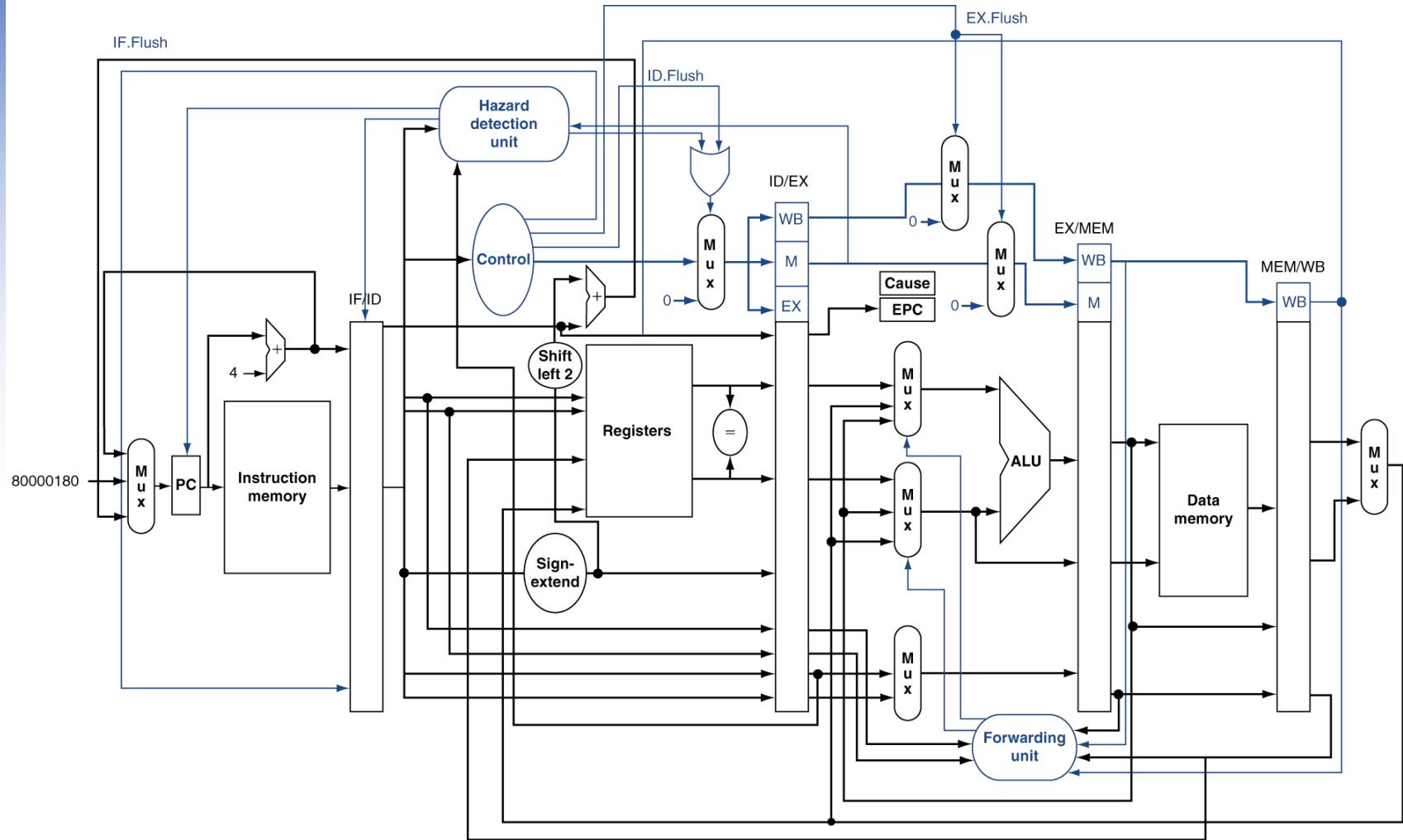
Handler Actions

- Read cause, and transfer to relevant handler
- Determine action required
- If restartable
 - Take corrective action
 - use EPC to return to program
- Otherwise
 - Terminate program
 - Report error using EPC, cause, ...

Exceptions in a Pipeline

- Another form of control hazard
- Consider overflow on add in EX stage
 - add \$1, \$2, \$1
 - Prevent \$1 from being clobbered
 - Complete previous instructions
 - Flush add and subsequent instructions
 - Set Cause and EPC register values
 - Transfer control to handler
- Similar to mispredicted branch
 - Use much of the same hardware

Pipeline with Exceptions



Exception Properties

- Restartable exceptions
 - Pipeline can flush the instruction
 - Handler executes, then returns to the instruction
 - Refetched and executed from scratch
- PC saved in EPC register
 - Identifies causing instruction
 - Actually $PC + 4$ is saved
 - Handler must adjust

Exception Example

- Exception on add in

40	sub	\$11,	\$2,	\$4
44	and	\$12,	\$2,	\$5
48	or	\$13,	\$2,	\$6
4C	add	\$1,	\$2,	\$1
50	slt	\$15,	\$6,	\$7
54	lw	\$16,	50(\$7)	

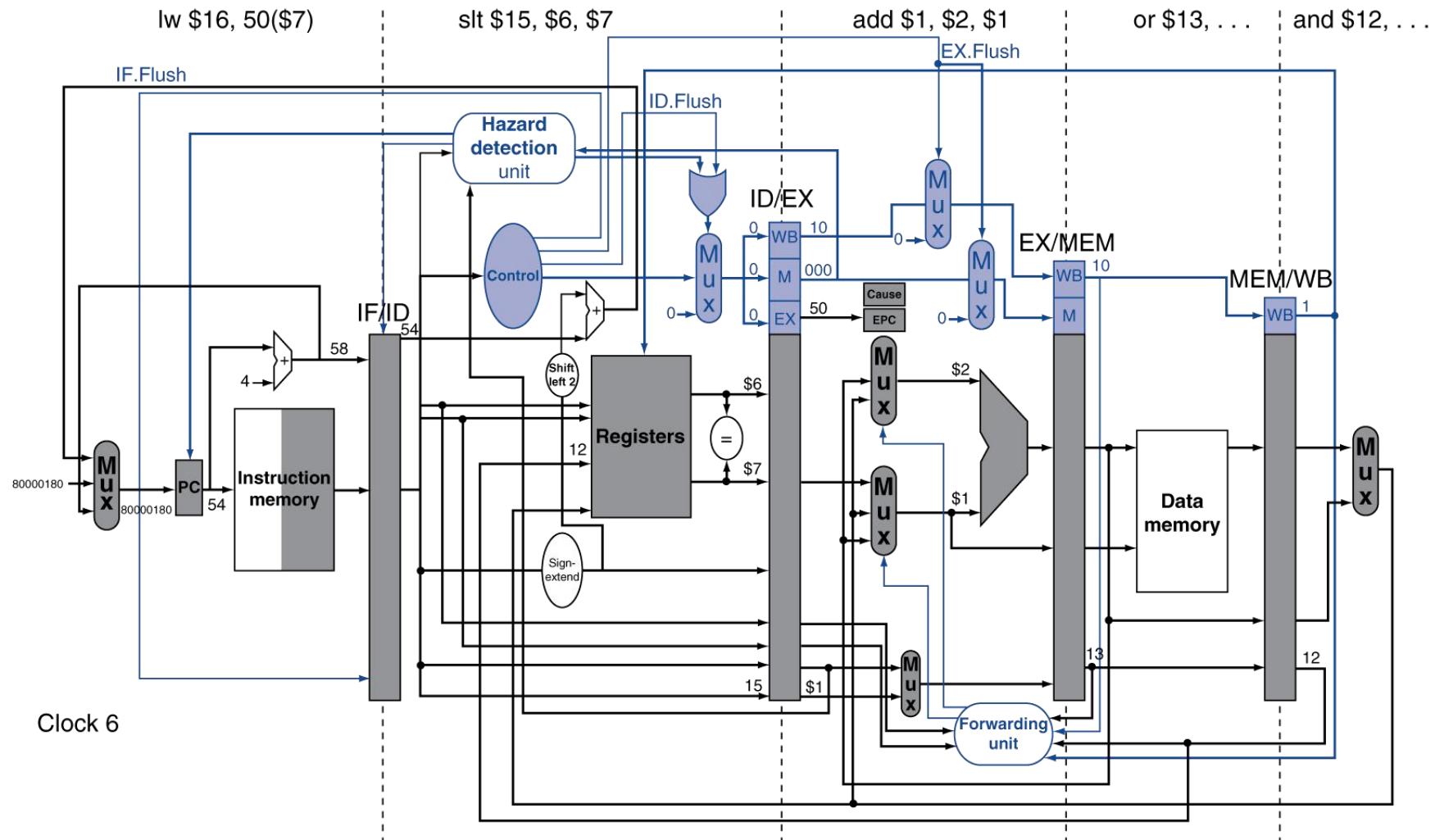
...

- Handler

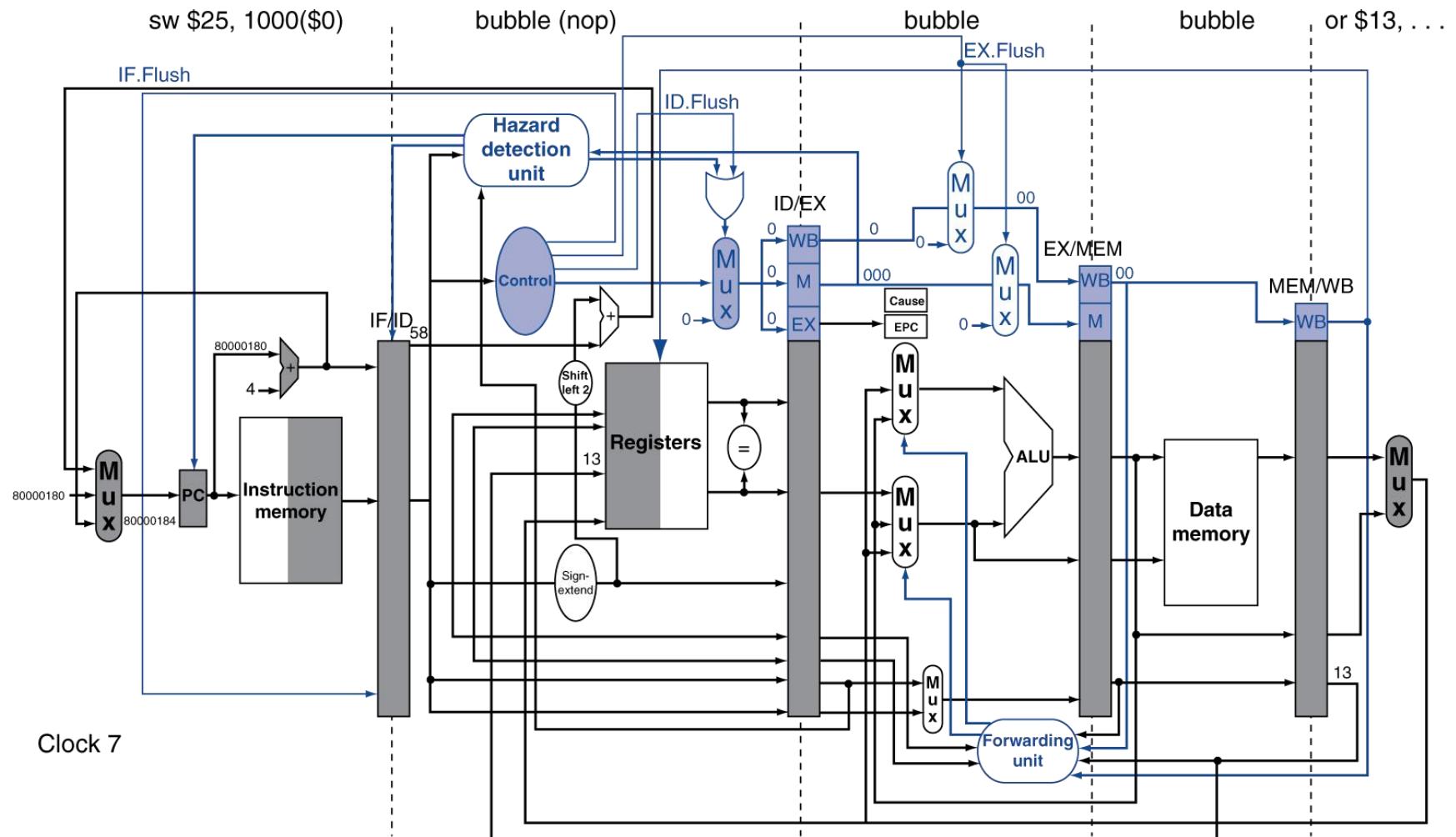
80000180	sw	\$25,	1000(\$0)
80000184	sw	\$26,	1004(\$0)

...

Exception Example



Exception Example



Multiple Exceptions

- Pipelining overlaps multiple instructions
 - Could have multiple exceptions at once
- Simple approach: deal with exception from earliest instruction
 - Flush subsequent instructions
 - “Precise” exceptions
- In complex pipelines
 - Multiple instructions issued per cycle
 - Out-of-order completion
 - Maintaining precise exceptions is difficult!

Imprecise Exceptions

- Just stop pipeline and save state
 - Including exception cause(s)
- Let the handler work out
 - Which instruction(s) had exceptions
 - Which to complete or flush
 - May require “manual” completion
- Simplifies hardware, but more complex handler software
- Not feasible for complex multiple-issue out-of-order pipelines

Three Generic Data Hazards

- **Read After Write (RAW)**

Instr_J tries to read operand before Instr_I writes it

For Now

I: add r1,r2,r3
J: sub r4, r1, r3

- Caused by a “Data Dependence” (in compiler nomenclature). This hazard results from an actual need for communication.

Three Generic Data Hazards

- Write After Read (WAR)

Instr_J writes operand before Instr_I reads it

```
I: sub r4, r1, r3  
J: add r1, r2, r3  
K: mul r6, r1, r7
```

- Called an “anti-dependence” by compiler writers.
This results from reuse of the name “r1”.
- Can't happen in MIPS 5 stage pipeline because:
 - All instructions take 5 stages, and
 - Reads are always in stage 2, and
 - Writes are always in stage 5

Three Generic Data Hazards

- **Write After Write (WAW)**

Instr_J writes operand before Instr_I writes it.

```
I: sub r1,r4,r3
J: add r1,r2,r3
K: mul r6,r1,r7
```

- Called an “**output dependence**” by compiler writers
This also results from the reuse of name “**r1**”.
- Can’t happen in MIPS 5 stage pipeline because:
 - All instructions take 5 stages, and
 - Writes are always in stage 5
- Will see WAR and WAW in later more complicated pipes

Review of Performance

Definitions

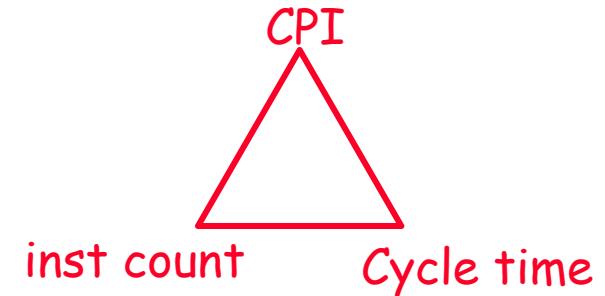
- Performance is in units of things per sec
 - bigger is better
- If we are primarily concerned with response time

$$\text{performance}(x) = \frac{1}{\text{execution_time}(x)}$$

" X is n times faster than Y" means

$$n = \frac{\text{Performance}(X)}{\text{Performance}(Y)} = \frac{\text{Execution_time}(Y)}{\text{Execution_time}(X)}$$

Computer Performance



$$\text{CPU time} = \frac{\text{Seconds}}{\text{Program}} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}$$

	Inst Count	CPI	Clock Rate
Program	X		
Compiler	X	(X)	
Inst. Set.	X	X	
Organization		X	X
Technology			X

Amdahl's Law

- Improving an aspect of a computer and expecting a proportional improvement in overall performance

$$T_{\text{improved}} = \frac{T_{\text{affected}}}{\text{improvement factor}} + T_{\text{unaffected}}$$

- Example: multiply accounts for 80s/100s
 - How much improvement in multiply performance to get 5× overall?

$$20 = \frac{80}{n} + 20$$

- Can't be done!

- Corollary: make the common case fast

Review of Memory Hierarchy

Principle of Locality

- Programs access a small proportion of their address space at any time
- Temporal locality
 - Items accessed recently are likely to be accessed again soon
 - e.g., instructions in a loop, induction variables
- Spatial locality
 - Items near those accessed recently are likely to be accessed soon
 - E.g., sequential instruction access, array data

Taking Advantage of Locality

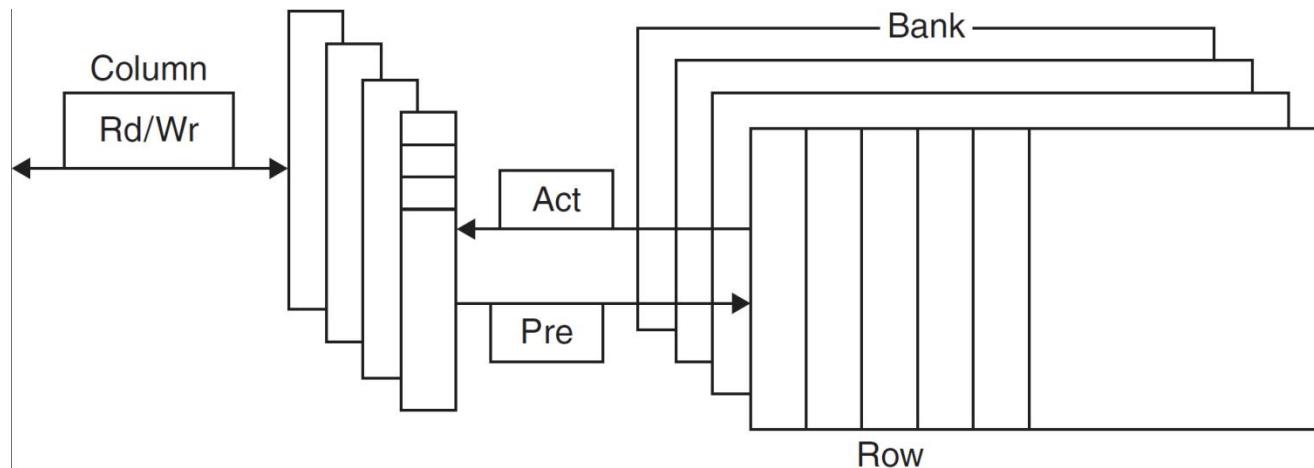
- Memory hierarchy
- Store everything on disk
- Copy recently accessed (and nearby) items from disk to smaller DRAM memory
 - Main memory
- Copy more recently accessed (and nearby) items from DRAM to smaller SRAM memory
 - Cache memory attached to CPU

Memory Technology

- Static RAM (SRAM)
 - 0.5ns – 2.5ns, \$2000 – \$5000 per GB
- Dynamic RAM (DRAM)
 - 50ns – 70ns, \$20 – \$75 per GB
- Magnetic disk
 - 5ms – 20ms, \$0.20 – \$2 per GB
- Ideal memory
 - Access time of SRAM
 - Capacity and cost/GB of disk

DRAM Technology

- Data stored as a charge in a capacitor
 - Single transistor used to access the charge
 - Must periodically be refreshed
 - Read contents and write back
 - Performed on a DRAM “row”

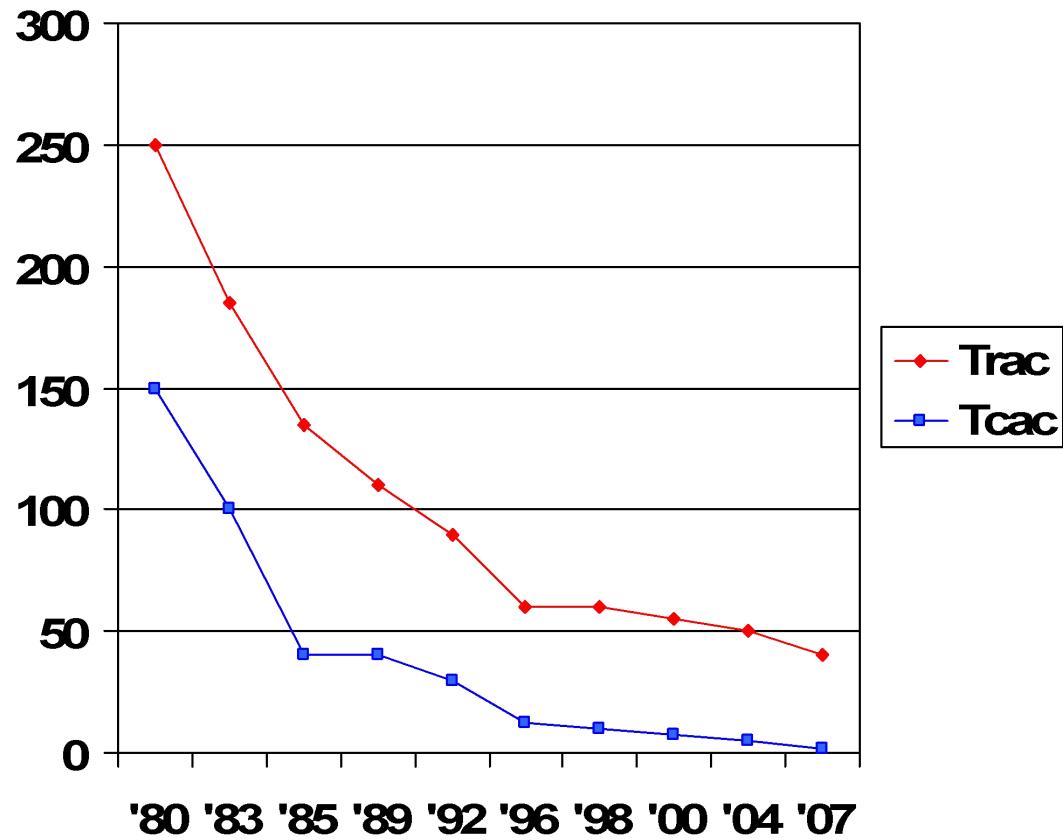


Advanced DRAM Organization

- Bits in a DRAM are organized as a rectangular array
 - DRAM accesses an entire row
 - Burst mode: supply successive words from a row with reduced latency
- Double data rate (DDR) DRAM
 - Transfer on rising and falling clock edges
- Quad data rate (QDR) DRAM
 - Separate DDR inputs and outputs

DRAM Generations

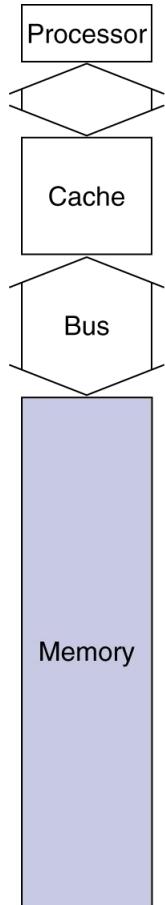
Year	Capacity	\$/GB
1980	64Kbit	\$1500000
1983	256Kbit	\$500000
1985	1Mbit	\$200000
1989	4Mbit	\$50000
1992	16Mbit	\$15000
1996	64Mbit	\$10000
1998	128Mbit	\$4000
2000	256Mbit	\$1000
2004	512Mbit	\$250
2007	1Gbit	\$50



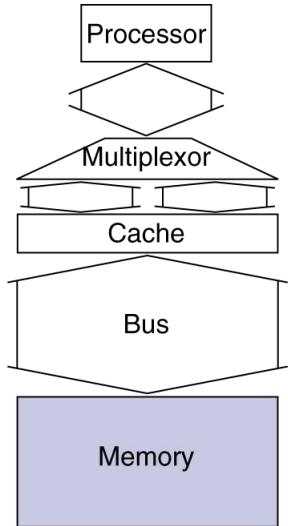
DRAM Performance Factors

- Row buffer
 - Allows several words to be read and refreshed in parallel
- Synchronous DRAM
 - Allows for consecutive accesses in bursts without needing to send each address
 - Improves bandwidth
- DRAM banking
 - Allows simultaneous access to multiple DRAMs
 - Improves bandwidth

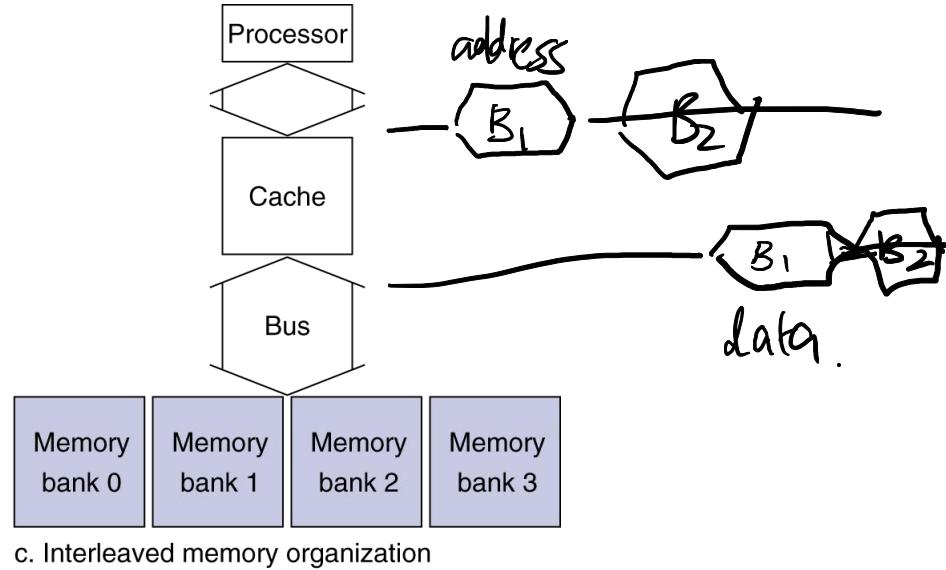
Increasing Memory Bandwidth



a. One-word-wide
memory organization



b. Wider memory organization



c. Interleaved memory organization

■ 4-word wide memory

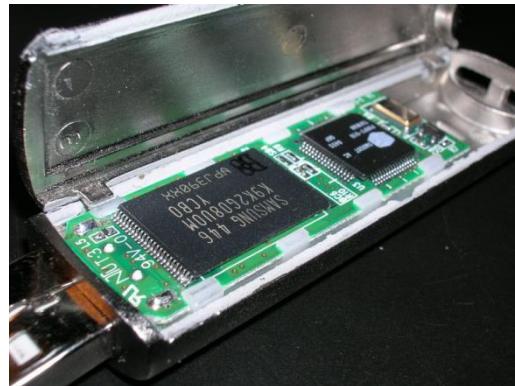
- Miss penalty = $1 + 15 + 1 = 17$ bus cycles
- Bandwidth = $16 \text{ bytes} / 17 \text{ cycles} = 0.94 \text{ B/cycle}$

■ 4-bank interleaved memory

- Miss penalty = $1 + 15 + 4 \times 1 = 20$ bus cycles
- Bandwidth = $16 \text{ bytes} / 20 \text{ cycles} = 0.8 \text{ B/cycle}$

Flash Storage

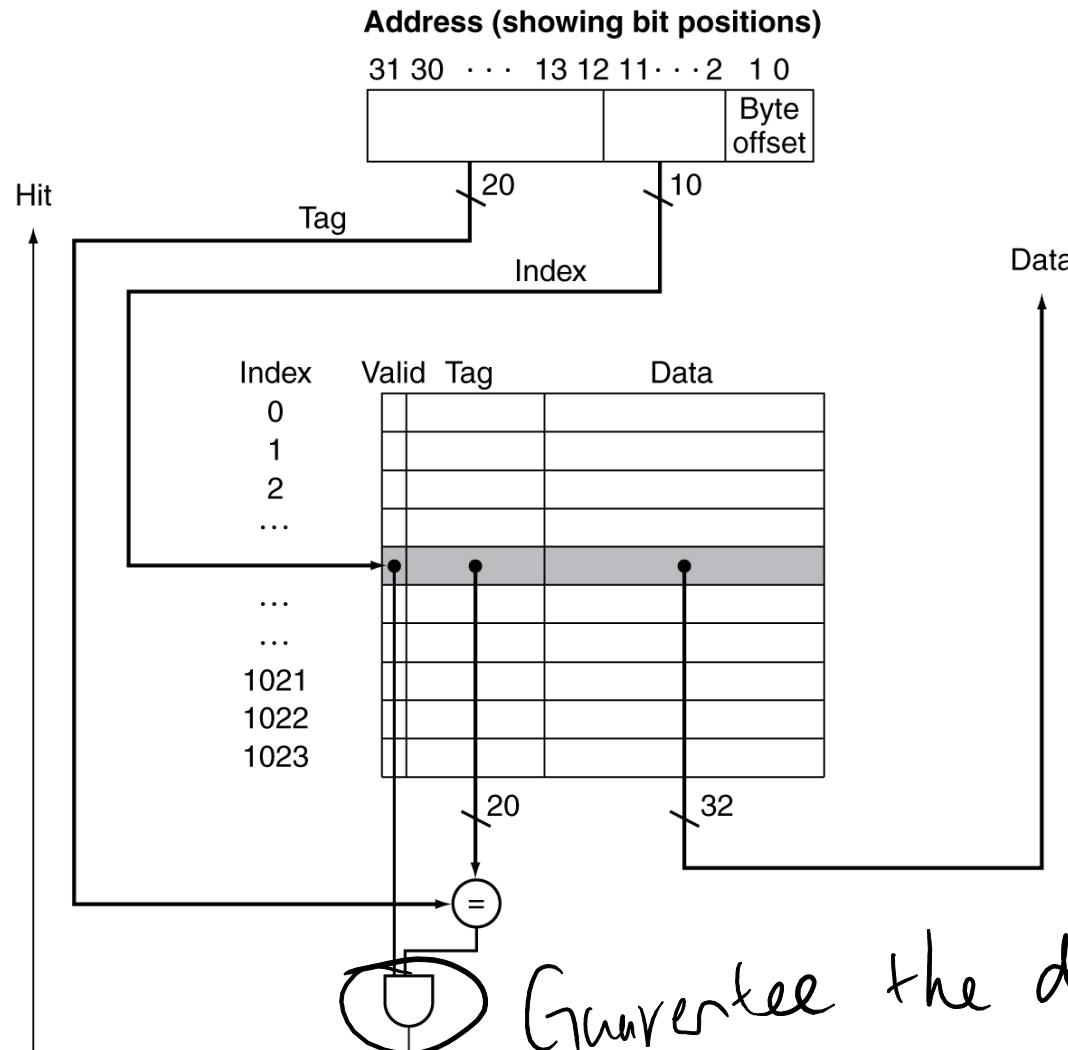
- Nonvolatile semiconductor storage
 - 100× – 1000× faster than disk
 - Smaller, lower power, more robust
 - But more \$/GB (between disk and DRAM)



Flash Types

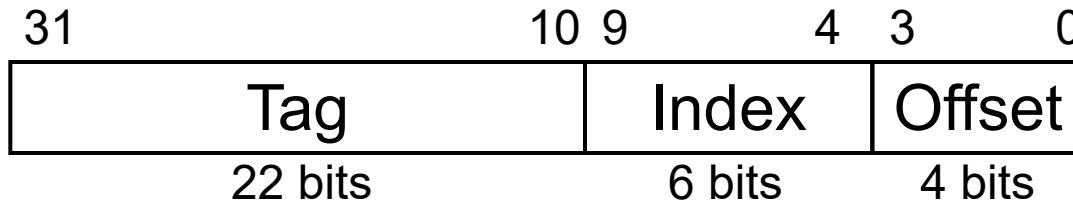
- NOR flash: bit cell like a NOR gate
 - Random read/write access
 - Used for instruction memory in embedded systems
- NAND flash: bit cell like a NAND gate
 - Denser (bits/area), but block-at-a-time access
 - Cheaper per GB
 - Used for USB keys, media storage, ...
- Flash bits wears out after 1000's of accesses
 - Not suitable for direct RAM or disk replacement
 - Wear leveling: remap data to less used blocks

Address Subdivision



Example: Larger Block Size

- 64 blocks, 16 bytes/block
 - To what block number does address 1200 map?
- Block address = $\lfloor 1200/16 \rfloor = 75$
- Block number = $75 \text{ modulo } 64 = 11$



Block Size Considerations

- Larger blocks should reduce miss rate
 - Due to spatial locality
- But in a fixed-sized cache
 - Larger blocks \Rightarrow fewer of them
 - More competition \Rightarrow increased miss rate
 - Larger blocks \Rightarrow pollution
- Larger miss penalty
 - Can override benefit of reduced miss rate
 - Early restart and critical-word-first can help

Cache Misses

- On cache hit, CPU proceeds normally
- On cache miss
 - Stall the CPU pipeline
 - Fetch block from next level of hierarchy
 - Instruction cache miss
 - Restart instruction fetch
 - Data cache miss
 - Complete data access

Write-Through

- On data-write hit, could just update the block in cache
 - But then cache and memory would be inconsistent
- Write through: also update memory
- But makes writes take longer
 - e.g., if base CPI = 1, 10% of instructions are stores, write to memory takes 100 cycles
 - Effective CPI = $1 + 0.1 \times 100 = 11$
- Solution: write buffer
 - Holds data waiting to be written to memory
 - CPU continues immediately
 - Only stalls on write if write buffer is already full

Write-Back

- Alternative: On data-write hit, just update the block in cache
 - Keep track of whether each block is dirty
- When a dirty block is replaced
 - Write it back to memory
 - Can use a write buffer to allow replacing block to be read first

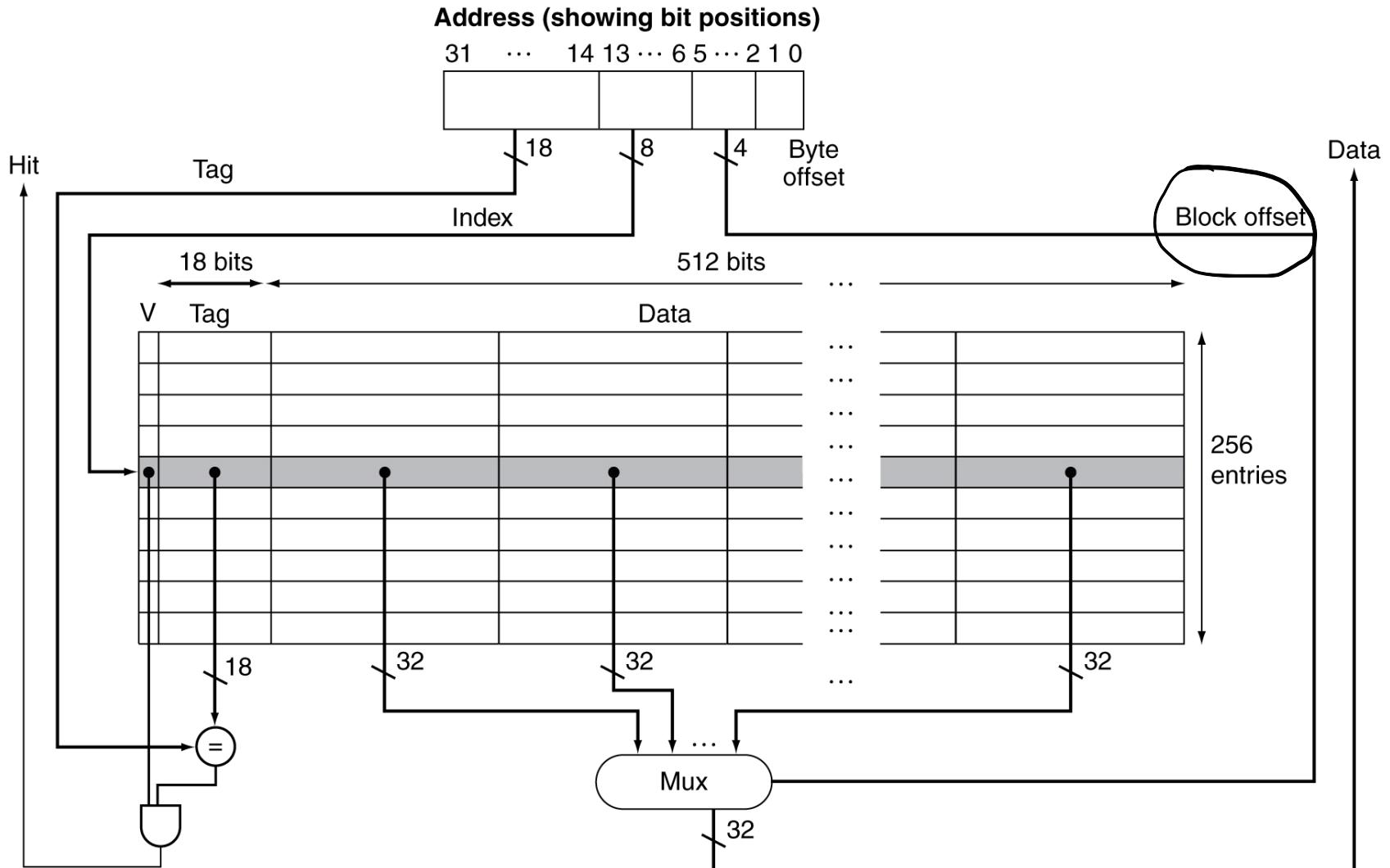
Write Allocation

- What should happen on a write miss?
- Alternatives for write-through
 - Allocate on miss: fetch the block
 - Write around: don't fetch the block
 - Since programs often write a whole block before reading it (e.g., initialization)
- For write-back
 - Usually fetch the block

Example: Intrinsity FastMATH

- Embedded MIPS processor
 - 12-stage pipeline
 - Instruction and data access on each cycle
- Split cache: separate I-cache and D-cache
 - Each 16KB: 256 blocks × 16 words/block
 - D-cache: write-through or write-back
- SPEC2000 miss rates
 - I-cache: 0.4%
 - D-cache: 11.4%
 - Weighted average: 3.2%

Example: Intrinsity FastMATH



Main Memory Supporting Caches

- Use DRAMs for main memory
 - Fixed width (e.g., 1 word)
 - Connected by fixed-width clocked bus
 - Bus clock is typically slower than CPU clock
- Example cache block read
 - 1 bus cycle for address transfer
 - 15 bus cycles per DRAM access
 - 1 bus cycle per data transfer
- For 4-word block, 1-word-wide DRAM
 - Miss penalty = $1 + 4 \times 15 + 4 \times 1 = 65$ bus cycles
 - Bandwidth = 16 bytes / 65 cycles = 0.25 B/cycle

Measuring Cache Performance

- Components of CPU time
 - Program execution cycles
 - Includes cache hit time
 - Memory stall cycles
 - Mainly from cache misses
- With simplifying assumptions:

Memory stall cycles

$$= \frac{\text{Memory accesses}}{\text{Program}} \times \text{Miss rate} \times \text{Miss penalty}$$

$$= \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Misses}}{\text{Instruction}} \times \text{Miss penalty}$$

Cache Performance Example

Given

- I-cache miss rate = 2%
- D-cache miss rate = 4%
- Miss penalty = 100 cycles
- Base CPI (ideal cache) = 2
- Load & stores are 36% of instructions

Miss cycles per instruction

- I-cache: $0.02 \times 100 = 2$
- D-cache: $0.36 \times 0.04 \times 100 = 1.44$

Actual CPI = $2 + 2 + 1.44 = 5.44$

- Ideal CPU is $5.44/2 = 2.72$ times faster

Average Access Time

- Hit time is also important for performance
- Average memory access time (AMAT)
 - $AMAT = \text{Hit time} + \text{Miss rate} \times \text{Miss penalty}$
- Example
 - CPU with 1ns clock, hit time = 1 cycle, miss penalty = 20 cycles, I-cache miss rate = 5%
 - $AMAT = 1 + 0.05 \times 20 = 2\text{ns}$
 - 2 cycles per instruction

Performance Summary

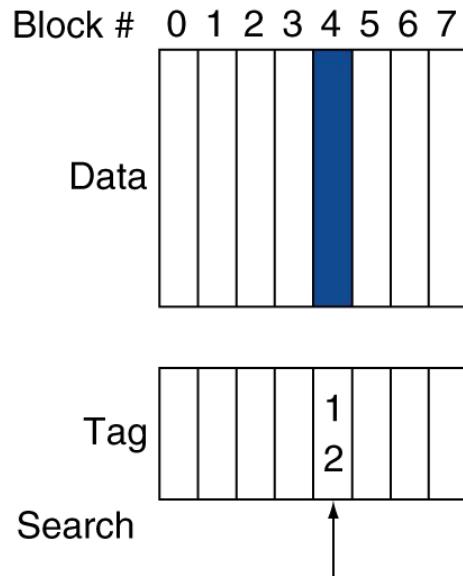
- When CPU performance increased
 - Miss penalty becomes more significant
- Decreasing base CPI
 - Greater proportion of time spent on memory stalls
- Increasing clock rate
 - Memory stalls account for more CPU cycles
- Can't neglect cache behavior when evaluating system performance

Associative Caches

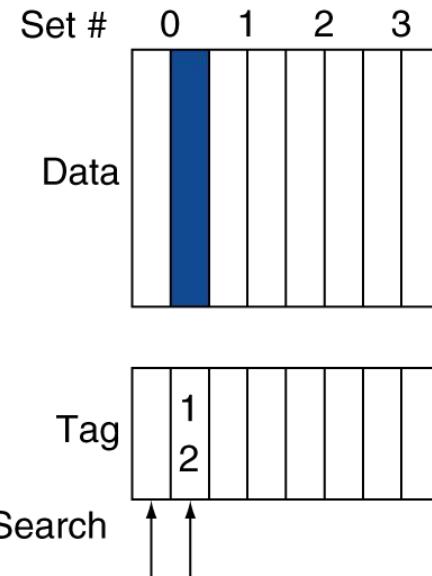
- Fully associative
 - Allow a given block to go in any cache entry
 - Requires all entries to be searched at once
 - Comparator per entry (expensive)
- n -way set associative
 - Each set contains n entries
 - Block number determines which set
 - (Block number) modulo (#Sets in cache)
 - Search all entries in a given set at once
 - n comparators (less expensive)

Associative Cache Example

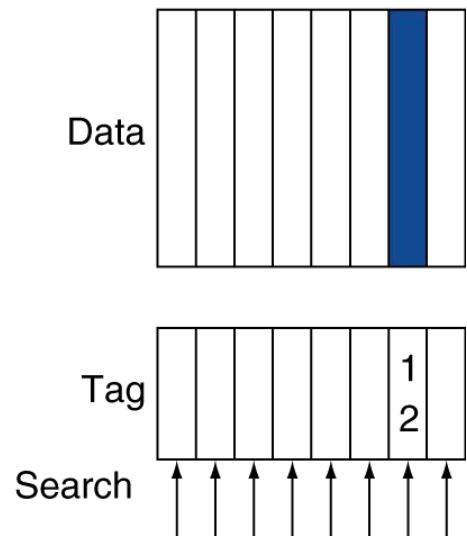
Direct mapped



Set associative



Fully associative



Spectrum of Associativity

- For a cache with 8 entries

**One-way set associative
(direct mapped)**

Block	Tag	Data
0		
1		
2		
3		
4		
5		
6		
7		

Two-way set associative

Set	Tag	Data	Tag	Data
0				
1				
2				
3				

Four-way set associative

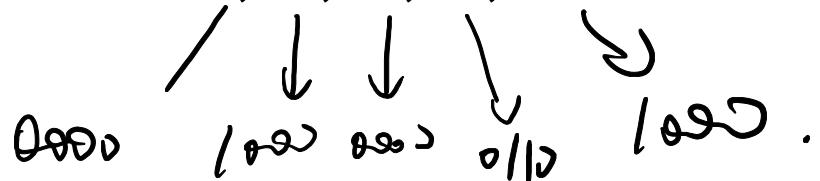
Set	Tag	Data	Tag	Data	Tag	Data	Tag	Data
0								
1								

Eight-way set associative (fully associative)

Tag	Data												

Associativity Example

- Compare 4-block caches
 - Direct mapped, 2-way set associative, fully associative
 - Block access sequence: 0, 8, 0, 6, 8
- Direct mapped



Block address	Cache index	Hit/miss	Cache content after access			
			0	1	2	3
0	0	miss	Mem[0]			
8	0	miss	Mem[8]			
0	0	miss	Mem[0]			
6	2	miss	Mem[0]		Mem[6]	
8	0	miss	Mem[8]		Mem[6]	

Associativity Example

■ 2-way set associative

Block address	Cache index	Hit/miss	Cache content after access		
			Set 0		Set 1
0	0	miss	Mem[0]		
8	0	miss	Mem[0]	Mem[8]	
0	0	hit	Mem[0]	Mem[8]	
6	0	miss	Mem[0]	Mem[6]	
8	0	miss	Mem[8]	Mem[6]	

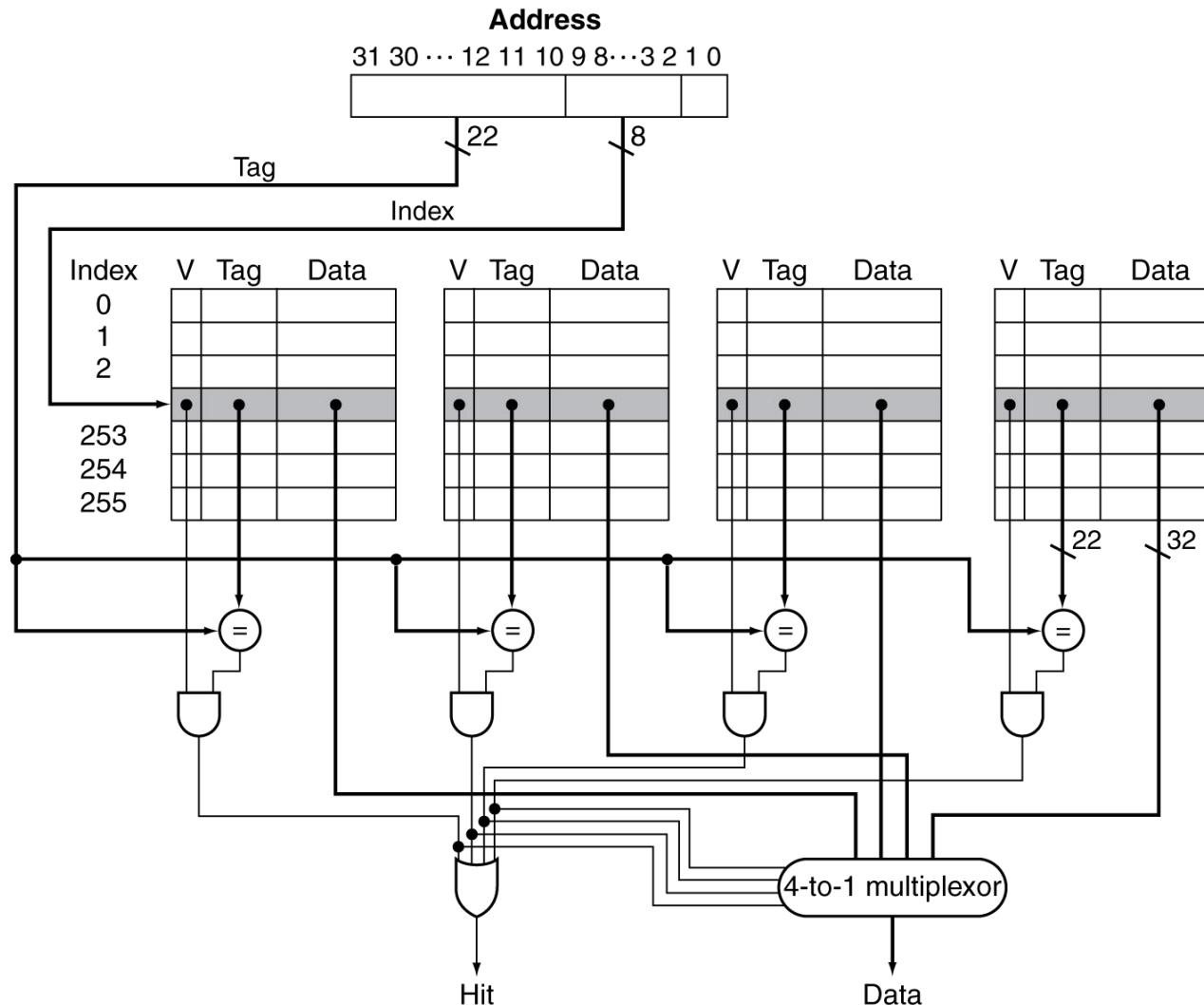
■ Fully associative

Block address		Hit/miss	Cache content after access			
0		miss	Mem[0]			
8		miss	Mem[0]	Mem[8]		
0		hit	Mem[0]	Mem[8]		
6		miss	Mem[0]	Mem[8]	Mem[6]	
8		hit	Mem[0]	Mem[8]	Mem[6]	

How Much Associativity

- Increased associativity decreases miss rate
 - But with diminishing returns
- Simulation of a system with 64KB D-cache, 16-word blocks, SPEC2000
 - 1-way: 10.3%
 - 2-way: 8.6%
 - 4-way: 8.3%
 - 8-way: 8.1%

Set Associative Cache Organization



Replacement Policy

- Direct mapped: no choice
- Set associative
 - Prefer non-valid entry, if there is one
 - Otherwise, choose among entries in the set
- Least-recently used (LRU)
 - Choose the one unused for the longest time
 - Simple for 2-way, manageable for 4-way, too hard beyond that
- Random
 - Gives approximately the same performance as LRU for high associativity

Multilevel Caches

- Primary cache attached to CPU
 - Small, but fast
- Level-2 cache services misses from primary cache
 - Larger, slower, but still faster than main memory
- Main memory services L-2 cache misses
- Some high-end systems include L-3 cache

Multilevel Cache Example

- Given
 - CPU base CPI = 1, clock rate = 4GHz
 - Miss rate/instruction = 2%
 - Main memory access time = 100ns
- With just primary cache
 - Miss penalty = $100\text{ns}/0.25\text{ns} = 400$ cycles
 - Effective CPI = $1 + 0.02 \times 400 = 9$

Example (cont.)

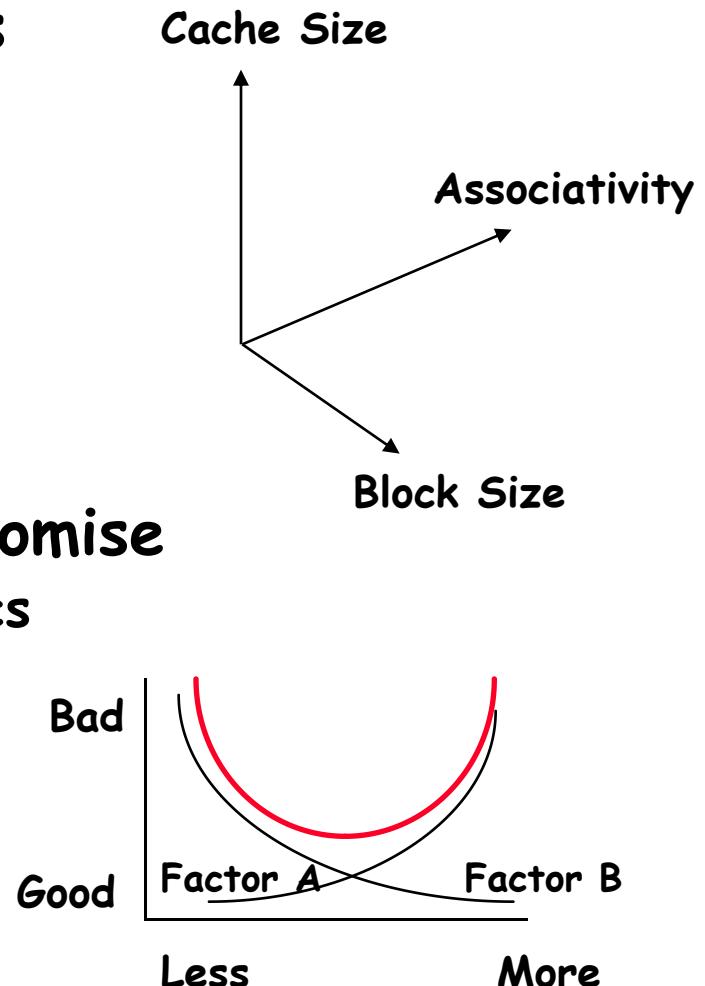
- Now add L-2 cache
 - Access time = 5ns
 - Global miss rate to main memory = 0.5%
- Primary miss with L-2 hit
 - Penalty = $5\text{ns}/0.25\text{ns} = 20 \text{ cycles}$
- Primary miss with L-2 miss
 - Extra penalty = 500 cycles
- CPI = $1 + 0.02 \times 20 + 0.005 \times 400 = 3.4$
- Performance ratio = $9/3.4 = 2.6$

Multilevel Cache Considerations

- Primary cache
 - Focus on minimal hit time
- L-2 cache
 - Focus on low miss rate to avoid main memory access
 - Hit time has less overall impact
- Results
 - L-1 cache usually smaller than a single cache
 - L-1 block size smaller than L-2 block size

The Cache Design Space

- Several interacting dimensions
 - cache size
 - block size
 - associativity
 - replacement policy
 - write-through vs write-back
- The optimal choice is a compromise
 - depends on access characteristics
 - » workload
 - » use (I-cache, D-cache, TLB)
 - depends on technology / cost
- Simplicity often wins

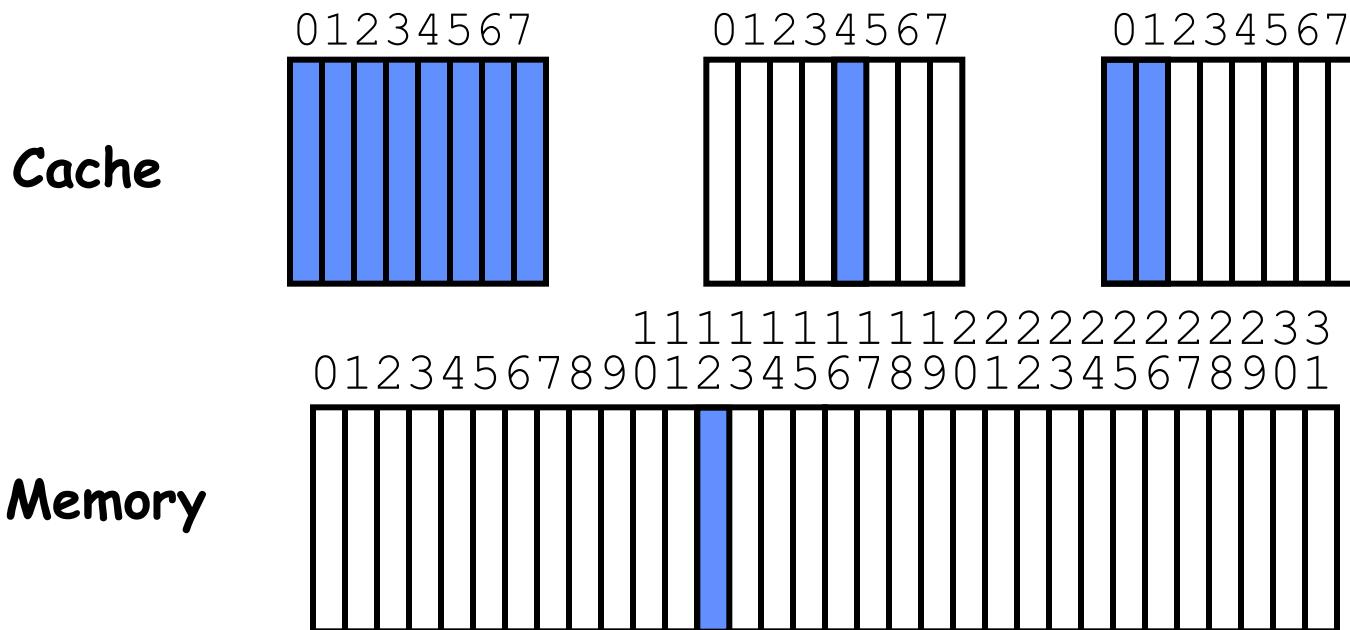


4 Questions for Memory Hierarchy

- Q1: Where can a block be placed in the upper level?
(Block placement)
- Q2: How is a block found if it is in the upper level?
(Block identification)
- Q3: Which block should be replaced on a miss?
(Block replacement)
- Q4: What happens on a write?
(Write strategy)

Q1: Where can a block be placed in the upper level?

- Block 12 placed in 8 block cache:
 - Fully associative, direct mapped, 2-way set associative
 - S.A. Mapping = Block Number Modulo Number
- | Sets | Full Mapped | Direct Mapped
$(12 \bmod 8) = 4$ | 2-Way Assoc
$(12 \bmod 4) = 0$ |
|------|-------------|-------------------------------------|-----------------------------------|
|------|-------------|-------------------------------------|-----------------------------------|



Q2: How is a block found if it is in the upper level?

- Tag on each block
 - No need to check index or block offset
- Increasing associativity shrinks index, → expands tag →



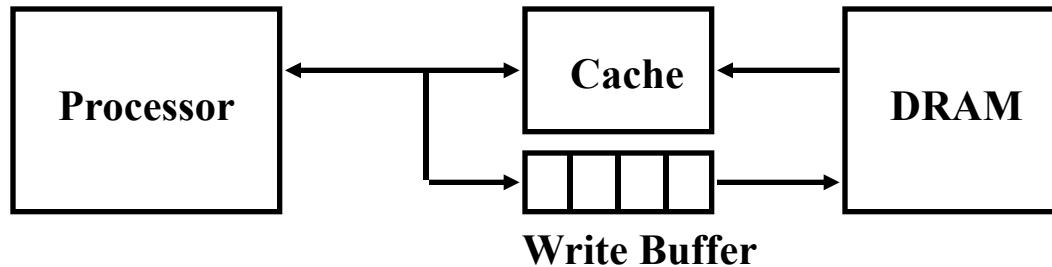
Q3: Which block should be replaced on a miss?

- **Easy for Direct Mapped**
- **Set Associative or Fully Associative:**
 - Random
 - LRU (Least Recently Used)
 - Other exotic techniques

Q4: What happens on a write?

- Write through—The information is written to both the block in the cache and to the block in the lower-level memory.
- Write back—The information is written only to the block in the cache. The modified cache block is written to main memory only when it is replaced.
 - is block clean or dirty?
- Pros and Cons of each?
 - WT: read misses cannot result in writes
 - WB: no repeated writes to same location
- WT always combined with write buffers so that don't wait for lower level memory

Write Buffer for Write Through



- A Write Buffer is needed between the Cache and Memory
 - Processor: writes data into the cache and the write buffer
 - Memory controller: write contents of the buffer to memory
- Write buffer is just a FIFO:
 - Typical number of entries: 4
 - Works fine if: $\text{Store frequency (w.r.t. time)} \ll 1 / \text{DRAM write cycle}$
- Memory system design:
 - $\text{Store frequency (w.r.t. time)} \rightarrow 1 / \text{DRAM write cycle}$
 - Write buffer saturation